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THE RELATIVE PERFORMANCE OF SCALABLE LOAD BALANCING ALGORITHMS IN LOOSELY-COUPLED DISTRIBUTED SYSTEMS

VOLII

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THE UNIVERSITY OF ASTON IN BIRMINGHAM

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APPENDIX A

LOAD BALANCING PROTOCOLS: MESSAGE TRACING

A.1 ROUTING TABLES

This section presents the message routing tables generated using Node0 as the initial starting point for traversing the network.

Processor	<u>-0</u> <u>PRIMA</u>	RY route	_	Processor1	PRIM	ARY route
dest.:0	via::node-1	distance::16 hops		dest.:0	via::node0	distance::1 hops
dest.:1	via::node1	distance::1 hops		dest.:1	via::node-1	distance::16 hops
dest.:2	via::node1	distance::2 hops		dest.:2	via::node2	distance::1 hops
dest.:3	via::node1	distance::3 hops		dest.:3	via::node2	distance::2 hops
dest.:4	via::node4	distance::I hops		dest.:4	via::node0	distance::2 hops
dest.:5	via::node1	distance::2 hops		dest.:5	via::node5	distance::1 hops
dest.:6	via::node4	distance::3 hops		dest.:6	via::node2	distance::2 hops
dest.:7	via::node4	distance::4 hops		dest.:7	via::node2	distance::3 hops
dest.:8	via::node4	distance::2 hops		dest.:8	via::node0	distance::3 hops
dest.:9	via::node1	distance::3 hops		dest.:9	via::node5	distance::2 hops
dest.:10	via::node1	distance::4 hops		dest.:10	via::node5	distance::3 hops
dest.:11	via::node4	distance::5 hops		dest.:11	via::node2	distance::4 hops
dest.:12	via::node4	distance::3 hops		dest.:12	via::node0	distance::4 hops
dest.:13	via::node1	distance::4 hops		dest.:13	via::node5	distance::3 hops
dest.:14	via::node1	distance::5 hops		dest.:14	via::node2	distance::4 hops
dest.:15	via::node4	distance::6 hops		dest.:15	via::node2	distance::5 hops
5	ROUTES	via this node		12	ROUTES	via this node ***

Table A.1 Message Routing Tables for Node0 and Node1

Processo	r2 PRIMAR	Y route		Processor3	PRIMARY	route
dest.:0	via::node1	distance::2 hops		dest.:0	via::node2	distance::3 hops
dest.:1	via::node1	distance::1 hops		dest.:1	via::node2	distance::2 hops
dest.:2	via::node-1	distance::16 hops		dest.:2	via::node2	distance::1 hops
dest.:3	via::node3	distance::1 hops		dest.:3	via::node-1	distance::16 hops
dest.:4	via::node1	distance::3 hops		dest.:4	via::node7	distance::4 hops
dest.:5	via::node6	distance::2 hops		dest.:5	via::node7	distance::3 hops
dest.:6	via::node6	distance::1 hops		dest.:6	via::node7	distance::2 hops
dest.:7	via::node3	distance::2 hops		dest.:7	via::node7	distance::1 hops
dest.:8	via::node1	distance::4 hops		dest.:8	via::node2	distance::5 hops
dest.:9	via::node1	distance::3 hops		dest.:9	via::node2	distance::4 hops
dest.:10	via::node6	distance::2 hops		dest.:10	via::node2	distance::3 hops
dest.:11	via::node6	distance::3 hops		dest,:11	via::node7	distance::2 hops
dest.:12	via::node6	distance::5 hops		dest.:12	via::node7	distance::6 hops
dest.:13	via::node6	distance::4 hops		dest.:13	via::node7	distance::5 hops
dest.:14	via::node6	distance::3 hops		dest.:14	via::node7	distance::4 hops
dest.:15	via::node3	distance::4 hops		dest.:15	via::node7	distance::3 hops
, , , , , , , , , , , , , , , , , , , ,			i,	And and the first of the second	***************************************	distance3 nops
11	ROUTES	via this node		6	ROUTES	via this node

Table A.2 Message Routing Tables for Node2 and Node3

Processor4	PRIMARY	route		Processor5	PRIMARY	route
dest.:0	via::node0	distance::I hops	1	dest.:0	via::node4	distance::2 hops
dest.:1	via::node5	distance::2 hops		dest.:1	via::node1	distance::1 hops
dest.:2	via::node0	distance::3 hops		dest.:2	via::node6	distance::2 hops
dest.:3	via::node5	distance::4 hops		dest.:3	via::node6	•
dest.:4	via::node-1	distance::16 hops		dest.:4	via::node4	distance::3 hops
dest.:5	via::node5	distance::1 hops		dest.:5	via::node-1	distance::1 hops
dest.:6	via::node5	distance::2 hops		dest.:6	via::node-1	distance::16 hops
dest.:7	via::node5	distance::3 hops		dest.:7	via::node6	distance::1 hops
dest.:8	via::node8	distance::1 hops		dest.:8		distance::2 hops
dest.:9	via::node8	distance::2 hops		dest.:9	via::node9	distance::2 hops
dest.:10	via::node8	distance::3 hops		dest.:10	via::node9	distance::1 hops
dest.:11	via::node8	distance::4 hops		dest.:11	via::node9	distance::2 hops
dest.:12	via::node8	distance::2 hops		dest.:12	via::node9	distance::3 hops
dest.:13	via::node8	distance::3 hops		dest.:12	via::node9	distance::3 hops
dest.:14	via::node5	distance::4 hops			via::node9	distance::2 hops
dest.:15	via::node5	distance::5 hops		dest.:14	via::node9	distance::3 hops
		distances nops	L	dest.:15	via::node9	distance::4 hops
9	ROUTES	via this node ***		14	ROUTES	via this node

Table A.3 Message Routing Tables for Node4 and Node5

Processo	<u>)r6 PRIMAR</u>	Y route	Processor	PRIMARY	
dest.:0	via::node5	distance::3 hops	dest.:0		route
dest.:I	via::node5	distance::2 hops	dest.:1	via::node3	distance::4 hops
dest.:2	via::node2	distance::1 hops	dest.:2	via::node3	distance::3 hops
dest.:3	via::node7	distance::2 hops	dest.:3	via::node3	distance::2 hops
dest.:4	via::node5	distance::2 hops	1 1=	via::node3	distance::1 hops
dest.:5	via::node5	distance::1 hops	dest.:4	via::node6	distance::3 hops
dest.:6	via::node-1	distance::16 hops	dest.:5	via::node6	distance::2 hops
dest.:7	via::node7	distance::1 hops	dest.:6	via::node6	distance::1 hops
dest.:8	via::node10	distance::3 hops	dest.:7	via::node-1	distance::16 hops
dest.:9	via::node10	distance::2 hops	dest.:8	via::nodel1	distance::4 hops
dest.:10	via::node10	distance::1 hops	dest.:9	via::node11	distance::3 hops
dest.:11	via::node10	distance::2 hops	dest.:10	via::node11	distance::2 hops
dest.:12	via::node10	distance::4 hops	dest.:11	via::nodel1	distance::1 hops
dest.:13	via::node10		dest.:12	via::node11	distance::5 hops
dest.:14	via::node10	distance::3 hops	dest.:13	via::node11	distance::4 hops
dest.:15	via::node10	distance::2 hops	dest.:14	via::node11	distance::3 hops
		distance::3 hops	dest.:15	via::node11	distance::2 hops
13	ROUTES	via this node ***	12	ROUTES	via this node ***

Table A.4 Message Routing Tables for Node6 and Node7

Processor	PRIMARY	<u>route</u>		Processor9	PRIMARY	route
dest.:0	via::node4	distance::2 hops	1	dest.:0	via::node8	distance::3 hops
dest.:1	via::node9	distance::3 hops		dest.:1	via::node5	distance::2 hops
dest.:2	via::node9	distance::4 hops		dest.:2	via::node10	distance::3 hops
dest.:3	via::node9	distance::5 hops		dest.:3	via::node10	distance::4 hops
dest.:4	via::node4	distance::1 hops		dest.:4	via::node8	distance::2 hops
dest.:5	via::node9	distance::2 hops		dest.:5	via::node5	distance::1 hops
dest.:6	via::node9	distance::3 hops		dest.:6	via::node10	distance::2 hops
dest.:7	via::node9	distance::4 hops		dest.:7	via::node10	distance::3 hops
dest.:8	via::node-1	distance::16 hops		dest.:8	via::node8	distance::1 hops
dest.:9	via::node9	distance::1 hops		dest.:9	via::node-1	distance::16 hops
dest.:10	via::node9	distance::2 hops		dest.:10	via::node10	distance::1 hops
dest.:11	via::node9	distance::3 hops		dest.:11	via::node10	distance::2 hops
dest.:12	via::node12	distance::1 hops		dest.:12	via::node13	distance::2 hops
dest.:13	via::node12	distance::2 hops		dest.:13	via::node13	distance::1 hops
dest.:14	via::node12	distance::3 hops		dest.:14	via::node13	distance::2 hops
dest.:15	via::node12	distance::4 hops		dest.:15	via::node13	distance::3 hops
12	ROUTES	via this node ***		23	ROUTES	via this node ***

Table A.5 Message Routing Tables for Node8 and Node9

Processo	r10 PRIMAR	RY route	Processor11	PRIMARY	route
dest.;0	via::node9	distance::4 hops	dest.:0	via::node10	distance::5 hops
dest.:1	via::node9	distance::3 hops	dest.:1	via::node10	distance::4 hops
dest.:2	via::node6	distance::2 hops	dest.:2	via::node10	distance::3 hops
dest.:3	via::node11	distance::3 hops	dest.:3	via::node7	distance::2 hops
dest.:4	via::node9	distance::3 hops	dest.:4	via::node10	distance::4 hops
dest.:5	via::node9	distance::2 hops	dest.:5	via::node10	distance::3 hops
dest.:6	via::node6	distance::1 hops	dest.:6	via::node10	distance::2 hops
dest.:7	via::node11	distance::2 hops	dest.:7	via::node7	distance::1 hops
dest.:8	via::node9	distance::2 hops	dest.:8	via::node10	distance::3 hops
dest.:9	via::node9	distance::1 hops	dest.:9	via::node10	distance::2 hops
dest.:10	via::node-1	distance::16 hops	dest.:10	via::node10	distance::1 hops
dest.:11	via::node11	distance::1 hops	dest.:11	via::node-1	distance::16 hops
dest.:12	via::node14	distance::3 hops	dest.:12	via::node15	distance::4 hops
dest.:13	via::node14	distance::2 hops	dest.:13	via::node15	distance::3 hops
dest.:14	via::node14	distance::1 hops	dest.:14	via::node15	distance::2 hops
dest.:15	via::node14	distance::2 hops	dest.:15	via::node15	distance::1 hops
23	ROUTES	via this node ***	16	ROUTES	via this node ***

Table A.6 Message Routing Tables for Node10 and Node11

	Processor1	2 PRIMARY	<u>route</u>		Processor13	<u>PRIMARY</u>	route		
			·	_					
	dest.:0	via::node8	distance::3 hops		dest.:0	via::node12	distance::4 hops		
	dest.:1	via::node13	distance::4 hops		dest.:1	via::node9	distance::3 hops		
	dest.:2	via::node13	distance::5 hops		dest.:2	via::node14	distance::4 hops		
	dest.:3	via::node8	distance::6 hops		dest.:3	via::node14	distance::5 hops		
	dest.:4	via::node8	distance::2 hops		dest.:4	via::node12	distance::3 hops		
	dest.:5	via::node13	distance::3 hops		dest.:5	via::node9	distance::2 hops		
	dest.:6	via::node13	distance::4 hops		dest.:6	via::node14	distance::3 hops		
	dest.:7	via::node8	distance::5 hops		dest.:7	via::node14	distance::4 hops		
	dest.:8	via::node8	distance::1 hops		dest.:8	via::node12	distance::2 hops		
	dest.:9	via::node13	distance::2 hops		dest.:9	via::node9	distance::1 hops		
	dest.:10	via::node13	distance::3 hops		dest.:10	via::node14	distance::2 hops		
	dest.:11	via::node13	distance::4 hops		dest.:11	via::node14	distance::3 hops		
	dest.:12	via::node-1	distance::16 hops		dest.:12	via::node12	distance::1 hops		
	dest.:13	via::node13	distance::1 hops		dest.:13	via::node-1	distance::16 hops		
	dest.:14	via::node13	distance::2 hops		dest.:14	via::node14	distance::1 hops		
	dest.:15	via::node13	distance::3 hops		dest.:15	via::node14	distance::2 hops		
•							The state of the s		
	7	ROUTES	via this node ***		20	ROUTES	via this node		

Table A.7 Message Routing Tables for Node12 and Node13

Processo	<u>r14 </u>	RY route		Processor15	PRIMA	RY route
dest.:0	via::node13	distance::5 hops	de	st.:0 via::	node14	distance::6 hops
dest.:1	via::node13	distance::4 hops	de	st.:1 via::	node11	distance::5 hops
dest.:2	via::node10	distance::3 hops	de	st.:2 via::	node11	distance::4 hops
dest.:3	via::node15	distance::4 hops	de	st.:3 via::	node11	distance::3 hops
dest.:4	via::node13	distance::4 hops	de	st.:4 via::	node14	distance::5 hops
dest.:5	via::node13	distance::3 hops	de	st.:5 via::	node14	distance::4 hops
dest.:6	via::node10	distance::2 hops	de	st.:6 via::	node11	distance::3 hops
dest.:7	via::node15	distance::3 hops	de	st.:7 via::	node11	distance::2 hops
dest.:8	via::node13	distance::3 hops	de	st.:8 via::	node14	distance::4 hops
dest.:9	via::node13	distance::2 hops	de	st.:9 via::	node14	distance::3 hops
dest.:10	via::node10	distance::1 hops	de	st.:10 via::	node11	distance::2 hops
dest.:11	via::node15	distance::2 hops	de	st.:11 via::	node11	distance::1 hops
dest.:12	via::node13	distance::2 hops	de	st.:12 via::	node14	distance::3 hops
dest.:13	via::node13	distance::1 hops	de	st.:13 via::	node14	distance::2 hops
dest.:14	via::node-1	distance::16 hops	de	st.:14 via::	node14	distance::1 hops
dest.:15	via::node15	distance::1 hops	de		node-1	distance::16 hops
			· • • • • • • • • • • • • • • • • • • •	iri er izanagini er erir iki e myakeeri er er azannea e er eraki eari	incontrol production and the control of the control	erioritation in the common to the state of t
18	ROUTES	via this node ***	7	ROL	JTES	via this node ***

Table A.8 Message Routing Tables for Node14 and Node15

A.2 ALGORITHMS

A.2.1 The Threshold Policy

The simulation model was run for the threshold policy with debugging information enabled. Statistics were obtained for the whole system and for each host every 500 seconds of simulated time. The following message dialogue was recorded shortly after the first phase of the simulation:

500.001.29152.90263.86202.2819 1.44 0.3155 7.3187 1.83

Node10 receives Packet N3-N10:: Node3 -overloaded- has asked node10 TO accept
Node10 GMT 500.040 TripT0.256 TIMER:WAITP_TOUT Etime: 500.296

Node10 receives Packet N3-N10:: node 10 RECEIVES process6396 [load = 2] FROM node 3

Node8 GMT 500.277 TripT0.260 TIMER:PRB TOUT Etime: 500.537

Node8 receives Packet N3-N8:: Node8::Time:500.582725 Node3's reply (prc6398) was 3

Node8 GMT 500.584 TripT0.254 TIMER:PRB_TOUT Etime: 500.838

Node10 receives Packet N8-N10:: Node8 -overloaded- has asked node10 TO accept Node10 GMT 500.612 TripT0.254 TIMER:WAITP TOUT Etime: 500.866

Node8 receives Packet N10-N8::

Node8::Time:500.735934 Node10's reply (prc6398) was -2 node 8 sends process6398 [load = 6] to node 10

Node10 receives Packet N8-N10:: node 10 RECEIVES process6398 [load = 3] FROM node 8 Node10 receives Packet N6-N10:: Node6 -overloaded- has asked node10 TO accept

Node3 probes node10 which subsequently accepts prc6396 from the latter. Node8 probes Node3 and node10 on behalf of prc6398 where the former rejects whilst the latter is willing to accept. A further request made by Node6 is rejected as Node10 is no longer underloaded.

Node8 receives Packet N9-N8:: node 8 receives Exit msg [prc:: 6392 at node:9]

Node8 receives Packet N11-N8:: node 8 receives Exit msg [prc:: 6393 at node:11]

Node8 GMT 502.410 TripT0.256 TIMER:PRB_TOUT Etime: 502.666

Node8 receives Packet N10-N8:: node 8 receives Exit msg [prc:: 6398 at node:10]

Node8 receives Packet N11-N8:: Node8::Time:502.585341 Node11's reply (prc6422) was 3

Node8 receives exit messages for processes 6392, 6393, and 6398 from nodes 9, 11, and 10 respectively. In addition, Node8 also probes Node11 and receives a rejection reply in return.

A.2.2 Communicating Set Threshold Policy

The simulation model was run using the same job stream as the Threshold policy but using the communicating set variation. The output produced by the model after the first period was:

```
SIMULATION no.1 STARTED ON Tue May 24 08:53:29 1994
  *** erand48 seeds are: 17757, 22851, and 10394 ***
                      16 Load value
  No. of Processors =
                                      = 0.80
  Total No. of jobs = 350000Ld Balancing Alg. = Threshold L/B Policy
  Threshold Level =
                    2.00 Probe Limit
  Convergence to < 2.00%
    TIME
             VAR. LOAD
                                DIFF. RTime %Conv
                                                         MIG.
                                                                Tx %JOB
   500.00
            1.4233 2.9512 4.0840 2.3243 2.43 0.3094 7.7293 1.83
N: 0 Mg:152.00
                  L: 3 Ri: 2.28
                                     Tx: 16 Imm: 121
                                                        Dih: 431
N: 1 Mg:171.00
                  L: 3 Ri: 2.45
                                                        Pth: 433
                                     Tx: 19 Imm: 145
N: 2 Mg:146.00
                  L: 3 Ri: 2.35
                                     Tx: 17 Imm: 160
                                                        Dih: 393
N: 3 Mg:148.00
                  L: 5 Rt: 2.50
                                     Tx: 14 Imm: 156
                                                        Dth: 401
N: 4 Mg:159.00
                  L: 3 Rt: 2.35
                                     Tx: 20 Imm: 163
                                                        Dth: 403
N: 5 Mg:156.00
                  L: 3 Rt: 2.17
                                     Tx: 23 Imm: 178
                                                        Dth: 388
N: 6 Mg:156.00
                  L: 3 Rt: 2.44
                                     Tx: 25 Imm: 155
                                                        Dth: 376
                  L: 2 Rt: 2.73
N: 7 Mg:167.00
                                     Tx: 25 Imm: 141
                                                        Dth: 400
N: 8 Mg:120.00
                  L: 4 Rt: 2.18
                                     Tx: 23 Imm: 163
                                                        Dth: 361
N: 9 Mg:152.00
                  L: 3 Rt: 2.12
                                     Tx: 36 Imm: 165
                                                        Dth: 389
N: 10 Mg:173.00
                  L: 3
                        Rt: 2.26
                                     Tx: 35 Imm: 176
                                                        Dth: 412
N: 11 Mg:150.00
                  L: 2 Rt: 2.30
                                     Tx: 27 Imm: 169
                                                        Dth: 377
N: 12 Mg:172.00
                  L: 3 Rt: 2.40
                                     Tx: 17 Imm: 152
                                                        Dth: 416
N: 13 Mg:142.00
                  L: 2 Rt: 2.59
                                     Tx: 34 Imm: 147
                                                        Dth: 370
N: 14 Mg:153.00
                  L: 3 Rt: 2.41
                                     Tx: 32 Imm: 159
                                                        Dth: 388
N: 15 Mg:158.00
                  L: 4 Rt: 2.59
                                     Tx: 24 Imm: 125
                                                        Dth: 410
```

Figure A.1 shows the average response time for each node after the first period of the simulation run. Nodes 5, 8, and 9 attained good response times, whilst node 7 produced the worse response time.

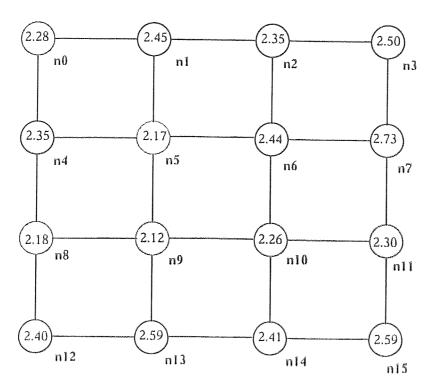


Figure A.1 Response Time After 500 Seconds using Cset Policy

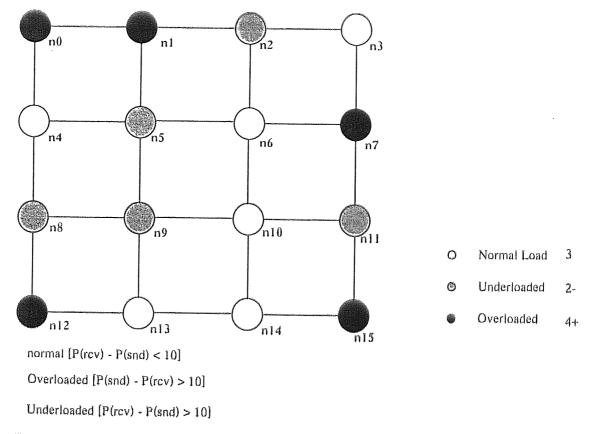


Figure A.2 Workload Distribution After 500 Seconds using Cset Policy

In Figure A.2 the load distribution in the network based on the ratio of processes sent and received is displayed.

The communicating set is built up as a by-product of load balancing events such as the receipt of or reply to a probe message. In this regard, node 10's view of the network in its overloaded state after the first period of the simulation run only identifies nodes 9 as an underloaded site for fruitful load sharing activities. In contrast, communicating set of Node 8 iis empty as it is unable to find an underloaded site.

```
N: 10 Mg:173.00 L: 3Rt: 2.26 Tx: 35 Imm: 176 Dth: 412 CommSet:[9->2]
```

```
N: 8 Mg:120.00 L: 4Rt: 2.18 Tx: 23 Imm: 163 Dth: 361 CommSet:[]
```

For ease of explanation and comprehension only the message dialogue of node 10 and node 8 is recorded. The dialogue is picked up three seconds before the end of the first simulation phase:

```
Node10 receives Packet N7-N10:: node 10 receives Exit msg [prc:: 6318 at node:7]

Node10 receives Packet N0-N10:: node 10 RECEIVES process6347 [load = 3] FROM node 0

Node10 receives Packet N0-N10:: Node0 -overloaded- has asked node10 TO accept

Node10 receives Packet N7-N10:: node 10 receives Exit msg [prc:: 6320 at node:7]

Node10 receives Packet N0-N10:: node 10 RECEIVES process6354 [load = 3] FROM node 0

Node10 receives Packet N14-N10:: Node14 -overloaded- has asked node10 TO accept
```

An exit message is received by node10 from node7 for one of its original processes (prc6318) followed by process 6347 from node0. As node10 is a member of the communicating set for node0 a further probe message is sent by the latter and the subsequent migration of prc6354 to node10. A further probe message by node14 was rejected as node10 is no longer underloaded.

```
Node8 receives Packet N6-N8::
                                node 8 receives Exit msg [prc:: 6341 at node:6]
Node8 receives Packet N12-N8::
                               Node12 -overloaded- has asked node8 TO accept
Node8 GMT 497.570
                         TripT0.252
                                      TIMER:WAITP_TOUT
                                                                  Etime: 497.822
Node8 receives Packet N12-N8::
                                node 8 RECEIVES process6356 [load = 2] FROM node 12
Node8 receives Packet N12-N8::
                               Node12 -overloaded- has asked node8 TO accept
Node8 GMT 497.644
                        TripT0.252
                                      TIMER:WAITP TOUT
                                                                  Etime: 497.896
Node8 receives Packet N12-N8::
                               node 8 RECEIVES process6359 [load = 3] FROM node 12
Node10 GMT 498.111
                        TripT0.254
                                      TIMER:PRB_TOUT Etime: 498,365
Node8 receives Packet N12-N8::
                               Node12 -overloaded- has asked node8 TO accept
```

Node8 receives consecutive probe messages from node12 resulting in the migration of processes 6356 and 6359 from the latter to the former. As node8 is a member of its communicating set a further probe message is sent. However, this will be rejected as Node8 is no longer underloaded.

Node10 receives Packet N7-N10:: Node10::Time:498.313249 Node7's reply (prc6364) was 3 Node10 GMT 498.314 TripT0.252 TIMER:PRB_TOUT Etime: 498.566 Node10 receives Packet N14-N10:: Node10::Time:498.365658 Node14's reply (prc6364) was 4 Node10 GMT 498.367 TripT0.252 TIMER:PRB_TOUT Etime: 498.619 Node10 receives Packet N14-N10:: Node14 -overloaded- has asked node10 TO accept

Node10 receives Packet N9-N10:; Node10::Time:498.424092 Node9's reply (prc6364) was -2 node 10 sends process6364 [load = 4] to node 9

As Node10 is now overloaded, its communicating set will be empty. That is, it would only contain information on overloaded sites that have sent it probe messages. Therefore, remote sites 7 and 14 are probed at random and both were found to be overloaded. A further probe message sent to node9 was successful resulting in the migration of prc6364.

Node8 receives Packet N9-N8:: Node9 -overloaded- has asked node8 TO accept Node8 receives Packet N14-N8:: node 8 receives Exit msg [prc:: 6344 at node:14] Node8 receives Packet N9-N8:: Node9 -overloaded- has asked node8 TO accept Node8 GMT 499.383 TripT0.256 TIMER:PRB TOUT Etime: 499.639 Node8 GMT 499,487 TripT0.256 TIMER:PRB TOUT Etime: 499.743 Node8 receives Packet N14-N8:: Node8::Time:499.588688 Node14's reply (prc6386) was 4 Node8 GMT 499.590 TripT0.252 TIMER:PRB TOUT Etime: 499.842 Node8 GMT 499.593 TripT0.254 TIMER:PRB_TOUT Etime: 499.847

Node8 receives Packet N4-N8:: Node8::Time:499.645715 Node4's reply (prc6386) was -2 node 8 sends process6386 [load = 5] to node 4

In its previous overloaded state node8 performed a successful migration (prc6344) to node14. Therefore, given its current state two probe messages are sent to node14 on behalf of the excess local processes. However, these requests are rejected and another two consecutive probe messages were sent to randomly selected hosts. Node4 had available capacity and process 6386 is migrated to it.

Node10 receives Packet N14-N10:: Node14 -overloaded- has asked node10 TO accept

Node10 receives Packet N0-N10:: Node0 -overloaded- has asked node10 TO accept
Node10 GMT 499.207 TripT0.258 TIMER:WAITP_TOUT Etime: 499.465
Node10 receives Packet N0-N10:: Node0 -overloaded- has asked node10 TO accept
Node10 GMT 499.456 TripT0.258 TIMER:WAITP_TOUT Etime: 499.714
Node10 receives Packet N0-N10:: node 10 RECEIVES process6384 [load = 3] FROM node 0

Node10 receives Packet N0-N10:: node 10 RECEIVES process6388 [load = 3] FROM node 0

Probe messages are sent by node14 and node0 to node10. However, at the time of node14's enquiry, the recipient was overloaded. Therefore, in the case of node0, the two requests made were accepted and process 6384 and 6388 migrated.

 Node8 GMT 499.659
 TripT0.252
 TIMER:PRB_TOUT
 Etime: 499.911

 Node8 receives Packet N14-N8::
 Node8::Time:499.710342
 Node14's reply (prc6391) was 4

 Node8 GMT 499.711
 TripT0.252
 TIMER:PRB_TOUT
 Etime: 499.963

Node8 receives Packet N0-N8:: Node8::Time:499.712351 Node0's reply (prc6392) was 3

Node8 GMT 499.713 TripT0.252 TIMER:PRB TOUT Etime: 499.965 Node8 receives Packet N4-N8:: Node8::Time:499.766364 Node4's reply (prc6393) was 3 Node8 GMT 499.767 TripT0.256 TIMER:PRB TOUT Etime: 500.023 Node8 receives Packet N4-N8:: Node8::Time:499.818773 Node4's reply (prc6391) was 3 Node8 GMT 499.820 TripT0.256 TIMER:PRB TOUT Etime: 500,076 Node8 receives Packet N4-N8:: Node8::Time:499.820782

Node8 receives Packet N4-N8:: Node8::Time:499.820782 Node4's reply (prc6392) was 3 Node8 GMT 499.822 TripT0.258 TIMER:PRB_TOUT Etime: 500.080

Node8 receives Packet N14-N8:: Node8::Time:499.923591 Node14's reply (prc6393) was 3 Node8 GMT 499.925 TripT0.258 TIMER:PRB TOUT Etime: 500.183

500.00 1.42332.95124.08402.3243 2.43 0.3094 7.7293 1.83

Node8 receives Packet N15-N8:: Node8::Time:500.026400 Node15's reply (prc6392) was 4
Node8 receives Packet N1-N8:: Node8::Time:500.027405 Node1's reply (prc6391) was 3
Node8 receives Packet N2-N8:: Node8::Time:500.179610 Node2's reply (prc6393) was 3
Node8 GMT 500.281 TripT0.254 TIMER:PRB_TOUT

Etime: 500.535

In this case node8 sends consecutive requests to the only member of its communicating set (node4) on behalf of processes 6391 to 6393, but receives in turn, consecutive rejections. This is also the case for requests sent to nodes 14, 0, 15, 1, and 2.

Node10 receives Packet N9-N10:: node 10 receives Exit msg [prc:: 6364 at node:9]

Node8 receives Packet N5-N8:: Node8::Time:500.484019 Node5's reply (prc6398) was 3

Node8 GMT 500.485 TripT0.256 TIMER:PRB_TOUT Etime: 500.741 Node10 receives Packet N6-N10:: Node6 -overloaded- has asked node10 TO accept

Node8 receives Packet N11-N8::

Node8::Time:500.637228 Node11's reply (prc6398) was -2 node 8 sends process6398 [load = 5] to node 11

Node8 receives Packet N4-N8:: node 8 receives Exit msg [prc:: 6386 at node:4]

Node8 receives Packet N15-N8:: Node15 -overloaded- has asked node8 TO accept

Both nodes 8 and 10 receive probe messages which they reject in turn as they are not in an underloaded state. However, node8 on its second attempt, found node11 willing to accept prc6398.

Node10 receives Packet N0-N10:: Node0 -overloaded- has asked node10 TO accept Node10 GMT 501.421 TripT0.258 TIMER:WAITP_TOUT Etime: 501.679 Node10 receives Packet N0-N10:: node 10 RECEIVES pr6412 [load = 2] FROM node 0

Node8 receives Packet N6-N8:: Node6 -overloaded- has asked node8 TO accept Node8 receives Packet N3-N8:: Node3 -overloaded- has asked node8 TO accept node8 receives Packet N11-N8:: node 8 receives Exit msg [prc:: 6398 at node:11]

Node10 receives Packet N4-N10:: Node4 -overloaded- has asked node10 TO accept Node10 GMT 502.188 TripT0.256 TIMER:WAITP_TOUT Etime: 502.444

Node10 receives Packet N4-N10:: node 10 RECEIVES prc6415 [load = 3] FROM node 4

Node8 GMT 502.383 TripT0.256 TIMER:PRB_TOUT Etime: 502.639

Node8 receives Packet N11-N8::

Node8::Time:502.585832 Node11's reply (prc6420) was -2

node 8 sends process6420 [load = 4] to node 11

As node10 is in an underloaded state the probes made to it by node0 and node4 were successful resulting in process migrations. Likewise, Node8 in an overloaded state also successfully probes the last node (nodell) to receive a process from it.

Node10 receives Packet N4-N10:: Node4 -overloaded- has asked node10 TO accept Node10 GMT 503.261 TripT0.256 TIMER:WAITP_TOUT Etime: 503.517 Node10 receives Packet N4-N10:: node 10 RECEIVES prc6432 [load = 3] FROM node 4

Node10 receives Packet N4-N10:: Node4 -overloaded- has asked node10 TO accept

Node8 receives Packet N11-N8:: node 8 receives Exit msg [prc:: 6420 at node:11]

Node8 receives Packet N4-N8:: Node4 -overloaded- has asked node8 TO accept

Node8 GMT 504.033 TripT0.252 TIMER:WAITP_TOUT Etime: 504.285

Node8 receives Packet N4-N8:: node 8 RECEIVES prc6437 [load = 3] FROM node 4

Node10 receives a further process from node4 but rejects its next request. Finally, node8 in in the underloaded state and accepts process 6437 from node4.

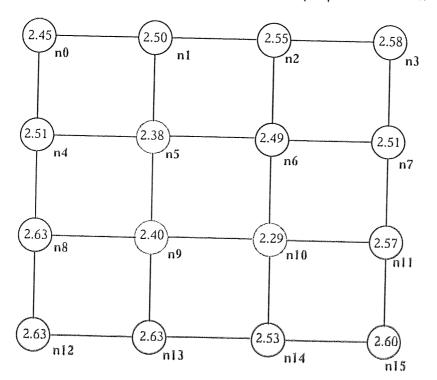


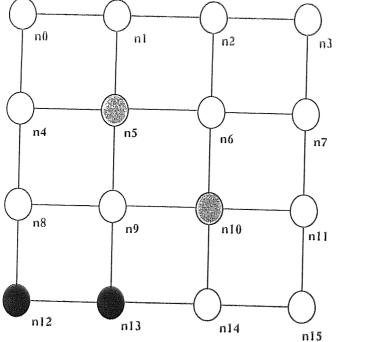
Figure A.3 Response Times After 4000 Seconds using Cset Policy

Figure A.3 shows the runtime performance after 4000 seconds of simulation time. Whilst the average response time has increased overall, the run-time performance of

nodes 3, 7, and 13 have improved. However, nodes 5, 9, and 10 continue to yield better than average performance despite the increase in response times.

In terms of their load state after 4000 seconds of simulated time, the following output was produced:

N:	0	Mg:1156.00	Ľ:	2	Rt: 2.45	Tx::	15	ľmm:1004	Dth: 3274
N:	1	Mg:1184.00	L:	3	Rt: 2.50	Tx::	19	Imm:1207	
N:	2	Mg:1200.00	L:	_	Rt: 2.55				Dth: 3257
		_		-		Tx::	18	ľmm:1243	Dth: 3200
N:		Mg:1154.00	L:	5	Ri: 2.58	Tx::	15	Imm:1178	Dth: 3179
N:	4	Mg:1197.00	L:	4	Rt: 2.51	Tx::	21	Imm:1242	Dth: 3213
N:		Mg:1270.00	L:	3	Rt: 2.38	Tx::	27	Imm:1329	Dth: 3272
N:	6	Mg:1217.00	L:	1	Rt: 2.49	Тх::	28	Imm:1229	Dth: 3208
N:	7	Mg:1255.00	L:	3	Rt: 2.51	Тх::	24	Imm:1254	Dth: 3197
N:	8	Mg:1211.00	L:	3	Rt: 2.63	Tx::	25	Imm:1186	Dth: 3184
N:	9	Mg:1286.00	Ľ:	3	Ri: 2.40	Tx::	38	Imm:1309	Dth: 3229
N:	10	Mg:1235.00	L:	2	Rt: 2.29	Tx::	38	Imm:1359	Dth: 3135
N:	11	Mg:1185.00	L:	ĺ	Rt: 2.57	Tx::	31	Imm:1212	Dth: 3092
N:	12	Mg:1186.00	L:	4	Rt: 2.63	Tx::	19	Imm:1135	Dth: 3236
N:	13	Mg:1240.00	L:	2	Rt: 2.63	Tx::	38		
		Mg:1232.00						Imm:1160	Dth: 3191
		-	L:	2	Ri: 2.53	Tx::	36	Imm:1193	Dth: 3226
N:	15	Mg:1213.00	L:	2	Rt: 2.60	Tx::	21	Imm:1181	Dth: 3215



- O Normal Load 3
- 9 Underloaded 2-
- Overloaded 4.

Overloaded [P(snd) - P(rcv) >> 50]

Underloaded [P(rcv) - P(snd) >>50]

Figure A.4 Workload Distribution After 4000 Seconds using Cset Policy

On the basis of the above table, nodes 3, 4, and 12 are overloaded and nodes 0, 2, 6, 10, 11, 13, 14, 15 underloaded. However, the ratio of the processes received and

transmitted identifies nodes 5 and 10 to be underloaded and nodes 12, and 13 to be overloaded. This is illustrated in Figure A.4.

A.2.3 Modified Communicating Set Threshold Policy

Given the weakness of thr previous communicating set implementation in handling clusters of process arrivals, the algorithm was modified such that random probes are selected until a reply is received for an initial probe to a member of the communicating set.

The output produced by the model after the first period was:

```
SIMULATION no.1 STARTED ON Tue May 24 08:53:29 1994
```

*** erand48 seeds are: 17757, 22851, and 10394 ***

No. of Processors = 16 Load value = 0.80

Total No. of jobs = 350000Ld Balancing Alg. = Threshold L/B Policy

Threshold Level = 2.00 Probe Limit = 3 Convergence to < 2.00%

TIME		VAR. LC	DAD	to 100	DIFF.	RTime	%Co	ηv	MIG.	Тх	%JO	В
500.00		1.3745 2.90	74 3.9	800	2.2420	3.00	0.324	1 7.4719	1.83			
N:	0	Mg:170.00	L:	5	Rt: 2.3	18 T	`x:	14	Imm: 126		Dth:	431
N:	1	Mg:177.00	L:	3	Rt: 2.1	9 T	x:	17	Imm: 160		Dth:	433
N:	2	Mg:156.00	L:	3	Rt: 2.2	r 8.	`x:	17	Imm: 149		Dth:	392
N:	3	Mg:169.00	L:	6	Rt: 2.2	5 T	`x:	11	Imm: 175		Dth:	400
N:	4	Mg:159.00	L:	2	Rt: 2.3	T 0	`x: 2	20	ľmm: 158		Dth:	403
N:	5	Mg:164.00	L:	3	Rt: 2.1	9 T	`x: 2	23	Imm: 181		Dth:	387
N:	6	Mg:159.00	L:	3	Rı: 2.2	5 T	'x: 2	23	Imm: 173		Dth:	377
N:	7	Mg:177.00	L:	4	Rt: 2.2	9 T	x: 2	20	Imm: 173		Dth:	398
N:	8	Mg:143.00	L:	6	Rt: 2.3	9 T	x: 1	19	Imm: 188		Dth:	361
N:	9	Mg:165.00	L:	3	Rt: 2.0	5 T	x: 3	32	Imm: 178		Dth:	388
N:	10	Mg:166.00	L:	1	Ri: 2.3	0 T	x: 3	38	Imm: 158		Dth:	412
N:	11	Mg:160.00	L:	3	Rt: 2.2	1 T	x: 2	24	Imm: 179		Dth:	378
N:	12	Mg:166.00	L:	4	Rt: 2.4	4 T	x: 1	6	Imm: 154		Dth:	415
N:	13	Mg:145.00	L:	3	Rt: 2.5	8 T	x: 3	13	Imm: 151		Dth:	370
N:	14	Mg:156.00	Ľ:	4	Rt: 2.3	8 T	x: 3	10	Imm: 166		Dth:	387
N:	15	Mg:161.00	L:	4	Rt: 2.5	0 Т	x: 2	23	Imm: 124		Dih:	408

Figure A.5 shows the average response time for each node after the first period of the simulation run. Nodes 1, 5, and 9 attained good response times, whilst nodes 12, 13, and 15 produced the worse response times.

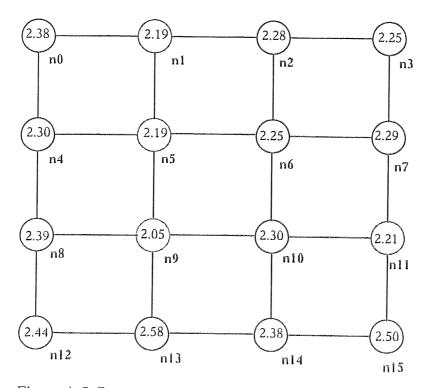
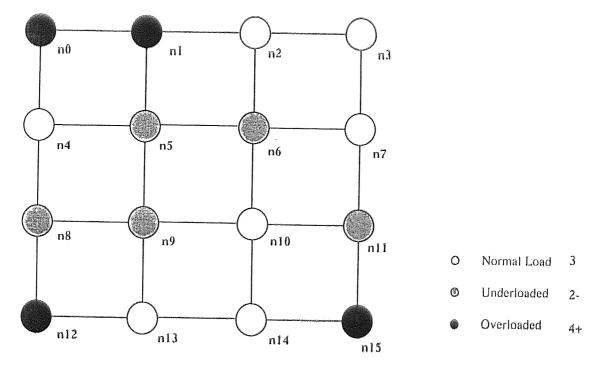


Figure A.5 Response Times After 500 Seconds using Modified Cset Implementation



normal [P(rcv) - P(snd) < 10]

Overloaded [P(snd) - P(rcv) > 10]

Underloaded [P(rcv) - P(snd) > 10]

Figure A.6 Workload After 500 Seconds (Modified Cset Implementation)

Figure A.6 shows the load distribution in the network based on the ratio of processes sent and received.

The following trace focussed on the dialogue for node θ in handling load balancing in cases where processes arrive in clusters.

Node8 receives Packet N9-N8:: Node9 -overloaded- has asked node8 TO accept Node8 receives Packet N12-N8:: Node12 -overloaded- has asked node8 TO accept Node8 GMT 499.386 TripT0.252 TIMER:PRB_TOUT Etime: 499.638 Node8 receives Packet N9-N8:: Node8::Time:499.489505 Node9's reply (prc6387) was 3 Node8 GMT 499.491 TripT0.254 TIMER:PRB TOUT Etime: 499.745 Node8 GMT 499,492 TripT0.256 TIMER:PRB TOUT Etime: 499,748

Node8 receives Packet N13-N8:: Node8::Time:499.648742 Node13's reply (prc6387) was 4

TIMER:PRB TOUT

Etime: 499.853

Four process arrivals at node8 resulted in the the random selection of four remote hosts to which probe messages were sent and their respective timers set. Node9 was the first to reply indicating that it was overloaded. Likewise, node13's reply was similar.

Node8 GMT 499.650 TripT0.258 TIMER:PRB TOUT Etime: 499.908 Node8 GMT 499.651 TripT0.256 TIMER:PRB TOUT Etime: 499.907

TripT0.256

Node8 GMT 499.597

Node8 receives Packet N11-N8:: Node8::Time:499.702155 Nodel1's reply (prc6391) was -1 node 8 sends process6391 [load = 7] to node 11

A further two processes arrive resulting in the random selection and probing of another two remote sites. The reply of Node11 indicates spare capacity, and process 6391 is migrated to it. In addition, Node11 becomes a member of the communicating set for Node8.

Node8 receives Packet N14-N8:: Node8::Time:499.767788 Node14's reply (prc6392) was 4 Node8 GMT 499,769 TripT0.256 TIMER:PRB_TOUT Etime: 500.025

Node8 receives Packet N6-N8:: Node8::Time:499.872607 Node6's reply (prc6393) was 3 Node8 GMT 499.874 TripT0.258 TIMER:PRB_TOUT Etime: 500.132

Node8 receives Packet N7-N8:: Node8::Time:499.925016 Node7's reply (prc6387) was 4

Node8 receives Packet N11-N8:: Node8::Time:499.926021 Nodell's reply (prc6392) was 3

A further two new processes are created, but node11 will be probed on behalf of the first process, removing its entry until a reply is received. The second process created must therefore select a site at random as the communicating set is now empty. The replies arriving from nodes 6, 7, and 11 indicated that they too were overloaded. Only the reply for the second process is outstanding.

In terms of their load state after 4000 seconds of simulated time, the following output was produced:

```
N: 0 Mg:1181.00
                  L: 3 Rt: 2.50
                                     Tx: 15
                                                  Imm: 989
                                                              Dth: 3274
N: 1 Mg:1251.00
                                     Tx: 18
                  L: 3 Rt: 2.43
                                                  Imm:1289
                                                              Dih: 3258
N: 2 Mg:1213.00
                  L: 2 R1: 2.42
                                     Tx: 17
                                                  Imm:1263
                                                              Dth: 3200
N: 3 Mg:1161.00
                  L: 4 Rt: 2.64
                                     Tx: 14
                                                  Imm:1183
                                                              Dth: 3179
N: 4 Mg:1229.00
                  L: 4 Rt: 2.45
                                     Tx: 20
                                                  Imm:1216
                                                              Dth: 3214
N: 5 Mg:1277.00
                  L: 2 Rt: 2.41
                                     Tx: 27
                                                  Imm:1284
                                                              Dth: 3272
N: 6 Mg:1227.00
                  L: 2 Rt: 2.38
                                     Tx: 26
                                                  Imm:1282
                                                              Dth: 3207
N: 7 Mg:1269.00
                  L: 3 Rt: 2.49
                                     Tx: 23
                                                              Dth: 3197
                                                  Imm:1279
N: 8 Mg:1279.00
                  L: 2 Rt: 2.47
                                     Tx: 22
                                                  Imm:1325
                                                              Dth: 3184
N: 9 Mg:1249.00
                  L: 3 Rt: 2.43
                                     Tx: 37
                                                  Imm:1317
                                                              Dth: 3229
N: 10 Mg:1223.00
                  L: 2 Rt: 2.36
                                     Tx: 38
                                                              Dth: 3136
                                                  Imm:1314
N: 11 Mg:1151.00
                  L: 2 Ri: 2.44
                                     Tx: 30
                                                  Imm:1237
                                                              Dth: 3091
N: 12 Mg:1161.00
                  L: 2 Rt: 2.52
                                     Tx: 18
                                                  Imm:1156
                                                              Dth: 3236
N: 13 Mg:1225.00
                  L: 3 RI: 2.66
                                     Tx: 38
                                                  Imm:1103
                                                              Dth: 3190
N: 14 Mg:1269.00
                  L: 3 Ri: 2.54
                                     Tx: 34
                                                  1mm:1210
                                                              Dth: 3225
N: 15 Mg:1250.00
                  L: 2 Ri: 2.63
                                     Tx: 21
                                                  Imm:1168
                                                              Dth: 3215
```

In terms of their load state after 6500 seconds of simulated time, the following output was produced:

650	0.0	00 1.5	5322	3.0	541	4.16	566 2.4	466	2.2	9	0.3023	7.843623.80		
N:	0	Mg:1776.0	00	L:	4	Rt:	2.49		Tx:	15	i	Imm:1651	Dth:	5197
N:	1	Mg:2006.0	00	L:	4	Rt:	2.47		Tx:	19)	Imm:2001	Dth:	5253
N:	2	Mg:2028.0	00	L:	4	Rt:	2.50		Tx:	18	3	Imm:1979	Dth:	5323
N:	3	Mg:1870.0	00	L:	2	Rt:	2.65		Tx:	14	ļ	Imm:1915	Dth:	5132
N:	4	Mg:1940.0	00	L:	1	Rt:	2.47		Tx:	20)	Imm:1997	Dth:	5161
N:	5	Mg:2024.0	00	L:	4	Rt:	2.42		Tx:	27	•	Imm:2099	Dth:	5268
N:	6	Mg:2029.0	00	L:	3	Rt:	2.46		Tx:	27	,	Imm:1969	Dih:	5272
N:	7	Mg:2002.0	00	L:	2	Rt:	2.52		Tx:	23	;	Imm:2040	Dth:	5192
N:	8	Mg:2015.0	00	L:	4	Rt:	2.52		Tx:	23		Imm:2041	Dth:	5194
N:	9	Mg:1963.0	00	L:	2	Rt:	2.45		Tx:	38		Imm:2089	Dth:	5178
N: 1	0	Mg:1963.0	00	Ľ:	3	Rt:	2.32		Tx:	37	'	Imm:2181	Dth:	5041
N: 1	1	Mg:1906.0	00	L:	2	Rı:	2.47		Tx:	31		Imm:1961	Dth:	5124
N: 1	2	Mg:1899.0	00	L:	3	Rı:	2.56		Tx:	19	•	Imm:1840	Dih:	5266
N: 1	13	Mg:1978.0	00	L:	2	Rt:	2.65		Tx:	38		Imm:1823	Dth:	5205
N: 1	4	Mg:2045.0	00	L:	2	Rt:	2.50		Tx:	34		lmm:1986	Dth:	5257
N: 1	15	Mg:1995.0	00	L:	1	Rı:	2.61		Tx:	21	1	lmm:1867	Dth:	5188

On the basis of the above table, nodes 0, 8, and 15 are overloaded and nodes 2, 11,12, and 13 underloaded. However, the ratio of processes received and transmitted identifies nodes 5, 9, and 10 to be underloaded and nodes 0, 8, 11, 13, and 14 to be overloaded. As thousands of processes are being considered, a host is considered to

be imbalanced if the difference in migration activity is significantly greater than fifty processes.

A.2.4 Global Average Neighbour (GsndNbor)

N: 15 Mg:168.00

The nearest-neighbour implementation of the sender-initiated global average algorithm attempts maintain a common and representative average load amongst its neighbours. Given the mesh topology, a by-product of its operation is the propogation of this average load throughout the network. The following extract represent the output from the model during the first period of the simulation run.

```
SIMULATION no.8 STARTED ON Thu May 19 17:28:10 1994
No. of Processors =
                    16
                          Load value
                                      = 0.80
Total No. of jobs = 350000 Ld Balancing Alg. = Global Avg Broadcast
Threshold Level =
                   2.00
                         Distance
Convergence to < 2.00%
TIME VAR. LOAD
                          DIFF. RTime %Conv
                                                    MIG.
                                                           Tx %JOB
                                .....
500.00 1.3489 2.7492 3.6720 2.0127 1.82 0.4036 5.1258 1.83
N: 0 Mg:171.00
                   L: 3 Rt: 2.20
                                      Tx: 13
                                                  Imm: 149
                                                               Dih: 431
N: 1 Mg:228.00
                   L: 4 Rt: 2.08
                                      Tx: 14
                                                  Imm: 189
                                                               Dth: 433
N: 2 Mg:200.00
                   L: 2 Rt: 2.05
                                      Tx: 11
                                                  Imm: 214
                                                               Dth: 393
N: 3 Mg:167.00
                   L: 2 Rt: 2.08
                                      Tx: 10
                                                  Imm: 165
                                                               Dth: 402
N: 4 Mg:214.00
                   L: 4 Rt: 2.11
                                      Tx: 13
                                                  Imm: 202
                                                               Dth: 402
N: 5 Mg:226.00
                   L: 2 Rt: 1.94
                                      Tx: 13
                                                  Imm: 243
                                                               Dth: 388
N: 6 Mg:196.00
                   L: 1 Rt: 1.93
                                      Tx: 13
                                                  Imm: 228
                                                               Dth: 377
N: 7 Mg:218.00
                   L: 3 Rt: 1.95
                                                               Dth: 400
                                      Tx: 11
                                                  Imm: 208
N: 8 Mg:179.00
                   L: 5 Rt: 2.00
                                                  Imm: 221
                                      Tx: 11
                                                               Dth: 361
N: 9 Mg:226.00
                  L: 2 Rt: 1.88
                                     Tx: 13
                                                  Imm: 237
                                                               Dth: 390
N: 10 Mg:243.00
                  L: 1 Rt: 2.01
                                                               Dth: 413
                                     Tx: 14
                                                  Imm: 221
N: 11 Mg:202.00
                                     Tx: 11
                   L: 2 Rt: 2.07
                                                  Imm: 227
                                                               Dth: 378
N: 12 Mg:190.00
                   L: 3 Rt: 2.07
                                     Tx: 11
                                                  Imm: 160
                                                               Dth: 417
N: 13 Mg:199.00
                  L: 3 Rt: 2.02
                                     Tx: 11
                                                  Imm: 219
                                                               Dth: 370
N: 14 Mg:202.00
                  L: 4 Rt: 2.17
                                     Tx: 13
                                                  Imm: 196
                                                               Dth: 388
```

L: 3 Rt: 2.27

Figure A.8 shows the average response time per host. The nodes offering the best average run time performance were nodes 5, 6, 7, and 9. The actual load state of the hosts after 500 seconds and displayed in Figure A.9, would indicate the overloaded hosts to be nodes 1, 3, 7, 8, and 13. Likewise, the underloaded hosts were nodes 4, 10, 11, and 12 with process queue length of two processes or less. Only node 10 had a response time that was well below the average.

Tx: 11

Imm: 150

Dth: 410

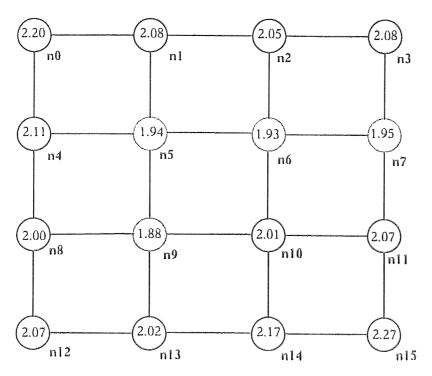


Figure A.8 Response Times After 500 Seconds using Global Average Policy

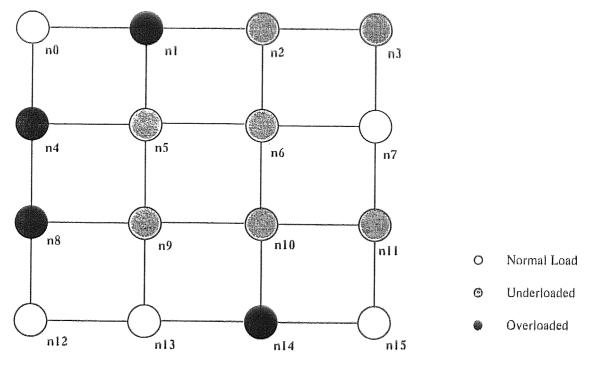
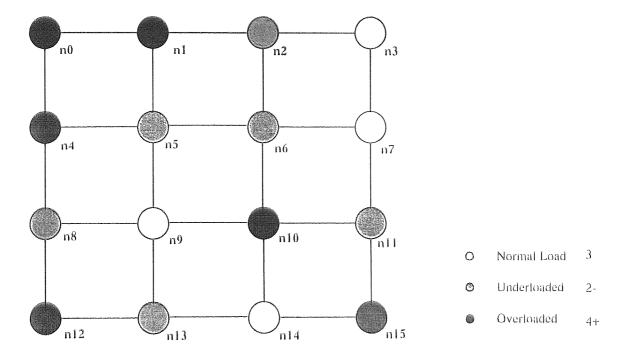


Figure A.9 Workload Distribution After 500 Seconds (Global Policy)



normal [P(rcv) - P(snd) < 10]

Overloaded [P(snd) - P(rcv) > 10]

Underloaded [P(rcv) - P(snd) > 10]

Figure A.10 Workload Distribution After 500 Seconds (Tx/Rx Ratio)

In Figure A.10, based on the ratio of processes sent and received over time, nodes θ , 1, 4, 10, 12, 15 would appear to be the overloaded nodes, whilst nodes 2, 5, 6, 8, 11, and 13 are underloaded.

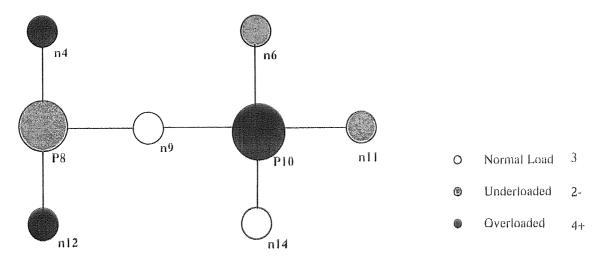


Figure A.11 Nearest-Neighbour Network Partition: Perspective of nodes 10 and 8

Consider the nodes 8 and 10 where the former is underloaded and the latter overloaded. Node8 has three neighbours, two of which are overloaded. The other

neighbour, node9 it has in common with node10. In contrast node 10 has four neighbours two of which are underloaded and two with normal loads. Figure A.11 shows the relationship between node 8 and 10 and their most immediate neighbours. The node (node9) common to both p8 and p10 is likely to be the point for the indirect propagation of the load average between the sites concerned.

The following trace focussed on the message dialogue of node10 and node8 just before and after the first period of the simulation run.

Node8 receives Packet N9-N8::

node 8 receives Exit msg [prc:: 6332 at node:9]

Node8 receives Packet N12-N8::

node 8 receives NewThreshold:2 [load:: 2] from node:12

Node8 receives Packet N12-N8::

Node12 -overloaded- has asked node8 TO accept

Node8 GMT 496.775

TripT0.252 TIMER: WAITP_TOUT Etime: 497.027

Node8 receives Packet N4-N8::

node 8 receives NewThreshold:2 [load:: 2] from node:4

Node8 receives Packet N12-N8::

node 8 receives Exit msg [prc:: 6326 at node:12]

WAITP TIMED OUT for Node8 at 497.030

Node8 receives exit messages for processes 6332 and 6326 from node9 and 12 respectively. It also received new load messages from node 12 and 4, setting the threshold level to two. As node8 is underloaded, a high load message from node12 resulted in the setting of the "await process" timer. Eventually, it times out as node12 may have chosen to migrate the process to one of its earlier respondents.

Node8 GMT 497.230

TripT0.252

TIMER:LOW_TOUT Etime: 497.482

Node10 GMT 497.386

TripT0.252

TIMER:LOW TOUT Etime: 497.638

LOWLD TIMED OUT for Node8 at 497,493

NewTh= 1

Node8 receives Packet N12-N8::

Node12 -overloaded- has asked node8 TO accept TripT0.252 TIMER:WAITP_TOUT

Etime: 497.756

LOWLD TIMED OUT for Node10 at 497.646

NewTh≈ 1

Node10 GMT 497.646

Node8 GMT 497.504

TripT0.252

TIMER:LOW_TOUT Etime: 497.898

Node10 receives Packet N11-N10:: Node11 -overloaded- has asked node10 TO accept Node10 GMT 497.677

TripT0.252

TIMER:WAITP_TOUT

Etime: 497.929

Node8 receives Packet N12-N8::

Node12 -overloaded- has asked node8 TO accept

WAITP TIMED OUT for Node8 at 497.767 WAITP TIMED OUT for Node10 at 497,938 Both nodes 8 and 10 are underloaded resulting in the setting of their respective "low load" timers. As node 8 had not received an enquiry from an overloaded host it eventually times out, assume the load threshold to be too high, and decrements and broadcast a new threshold value of one. As node12 is an immediate neighbour, it now becomes overloaded and sends an enquiry to all its neighbours (including node8). A similar pattern of behaviour is exhibited by node10. However, both sites were not selected as recipient sites and subsequently timed out.

Node8 GMT 497.987 TripT0.252 TIMER:LOW_TOUT Etime: 498.239 Node8 receives Packet N12-N8:: node 8 receives NewThreshold:2 [load:: 1] from node:12

Node8 receives Packet N12-N8:: Node12 -overloaded- has asked node8 TO accept Node8 GMT 498.019 TripT0.252 TIMER:WAITP_TOUT Etime: 498.271 Node8 GMT 498.240 TripT0.252 TIMER:LOW_TOUT Etime: 498.492

Node8 receives Packet N12-N8:: Node12 -overloaded- has asked node8 TO accept Node8 GMT 498.281 TripT0.252 TIMER:WAITP_TOUT Etime: 498.533

Node8 receives Packet N12-N8:: node 8 RECEIVES process6365 [load = 3] FROM node 12

As node 8 is still underloaded, it sets "low load" timer, but receives a new threshold value of two from node12. Consecutive enquiries by node12 resulted in the migration of process 6365 to node8.

Node10 receives Packet N14-N10:: Node14 -overloaded- has asked node10 TO accept Node8 receives Packet N4-N8:: Node4 -overloaded- has asked node8 TO accept

WAITP TIMED OUT for Node8 at 498.568

Node8 receives Packet N4-N8:: node 8 receives Exit msg [prc:: 6344 at node:4]

Node8 receives Packet N9-N8:: Node9 -overloaded- has asked node8 TO accept
Node8 GMT 498.721 TripT0.252 TIMER:WAITP_TOUT Etime: 498.973

Node10 receives Packet N9-N10:: Node9 -overloaded- has asked node10 TO accept Node10 receives Packet N14-N10:: Node14 -overloaded- has asked node10 TO accept

Node8 receives Packet N9-N8:: Node9 -overloaded- has asked node8 TO accept

WAITP TIMED OUT for Node8 at 498.976

Node 10 has received enquiries from nodes 14 and 9. These are rejected as there is no spare capacity on node10. Node 8 also received enquiries in sequence from nodes 4 and 9, but only the latter invoked a "wait for process" timer. However, node 9 decides not to send the process to node8 and it therefore times out.

Node10 receives Packet N14-N10:: node 10 receives NewThreshold:2 [load:: 2] from node:14

Node10 receives Packet N9-N10:: Node9 -overloaded- has asked node10 TO accept Node10 GMT 498.962 TripT0.252 TIMER:WAITP_TOUT Etime: 499.214 Node10 receives Packet N9-N10::

node 10 RECEIVES process6370 [load = 3] FROM node 9

Node8 GMT 499.027 TripT0.252 TIMER:LOW_TOUT Etime: 499.279 LOWLD TIMED OUT for Node8 at 499.329 NewTh= 2

Node10 GMT 499.430 TripT0.252 TIMER:LOW_TOUT Etime: 499.682 LOWLD TIMED OUT for Node10 at 499.732 NewTh= 2

Node 10 receives a new threshold value of two from node14. An enquiry is received from node9 resulting in the subsequent migration of process 6370 to node10. Nodes 8 and 10 become underloaded, set their respective "low load" timer, and both time out as a result of the absence of any enquiries by overloaded sites. The current threshold of three is assumed to be too high and is therefore decremented and broadcast to neighbouring sites.

Node8 GMT 499.787 TripT0.252 TIMER:HIGH_TOUT Etime: 500.039

Node10 receives Packet N9-N10::

node 10 receives NewThreshold:2 [load:: 2] from node:9

Node8 receives Packet N9-N8::

node 8 receives NewThreshold:2 [load:: 5] from node:9

Node8 receives Packet N4-N8:: (4) Node8 receives Packet N12-N8::(2)

Node8 receives Packet N9-N8:: (2)

Node10 GMT 499.936 TripT0.252 TIMER:LOW_TOUT Etime: 500.188

Node8 is overloaded, sets the "high load" timer and broadcasts an implicit request for assistance to its neighbours. Both node8 and 9 receive a new load message from node 9 which results in the cancellation of any load timers. The replies received by node8 from nodes 4, 12, and9 to its "high load" enquiry would indicate that both 12 and 9 had spare capacity. However, in light of the earlier change in the threshold all three respondents are considered as being overloaded.

500.00 1.3489 2.7492 3.6720 2.0127 1.82 0.4036 5.1258 1.83

Node8 GMT 500.046 TripT0.252 TIMER:HIGH_TOUT Etime: 500.298

Node8 receives Packet N4-N8:: (4) Node8 receives Packet N12-N8:: (4) Node8 receives Packet N9-N8:: (3)

LOWLD TIMED OUT for Node10 at 500.200 NewTh= 1

Node10 receives Packet N14-N10:: Node14 -overloaded- has asked node10 TO accept Node10 GMT 500.241 TripT0.252 TIMER:WAITP_TOUT Etime: 500.493

HIGHLD TIMED OUT for Node8 at 500.304 NewTh= 3
Node8 GMT 500.307 TripT0.252 TIMER:HIGH_TOUT Etime: 500.559

Node10 receives Packet N14-N10:: node 10 RECEIVES process6387 [load = 2] FROM node 14

Node8 receives Packet N12-N8::(2) node 8 sends process6398 [load = 6] to node 12 Node8 receives Packet N9-N8::(2) Node8 receives Packet N4-N8::(3)

As node8 is overloaded a "high load" message is broadcast to its neighbours. However, the replies indicate that nodes 4, 12, and 9 are overloaded. Therefore, node 8 times out, assumes the threshold to be too low, and increments and broadcast a new threshold value of three. When the algorithm is next executed, node8 is still overloaded, but two of its neighbours (nodes 12 and 9) are underloaded. Process 6387 is migrated to the first respondent, node12. In contrast, an underloaded node 10 assume the threshold to be too high, decrements and broadcast a threshold value of one which results in the migration of process 6387 from node14.

Node8 GMT 500.578 TripT0.252 TIMER:HIGH_TOUT Etime: 500.830

Node10 receives Packet N14-N10:: Node14 -overloaded- has asked node10 TO accept Node10 receives Packet N6-N10:: Node6 -overloaded- has asked node10 TO accept Node8 receives Packet N4-N8:: Node4 -overloaded- has asked node8 TO accept

Node8 receives Packet N12-N8::(3) node 8 sends process6392 [load = 5] to node 12

Node8 receives Packet N9-N8:: (2) Node8 receives Packet N4-N8::(3)

Node8 receives Packet N4-N8:: node 8 receives NewThreshold:2 [load:: 4] from node:4 Node8 GMT 500.851 TripT0.252 TIMER:HIGH TOUT Etime: 501.103

Node10 receives Packet N6-N10:: node 10 receives NewThreshold:2 [load:: 2] from node:6

Node 8 is overloaded and broadcast to its neighbours. The response by all neighbours indicates there load to be below the threshold, but only the first respondent is selected as recipient for the excess process load. The subsequent arrival of a new load message from neighbouring node4 would indicate that it has been unsuccessful in receiving remote processes, assumed the threshold to be too high and therefore broadcast a lower threshold value. In a similar manner, Node 10 received enquiries from overloaded sites 14 and 6. As these were rejected, a new load message duly arrived from node 6.

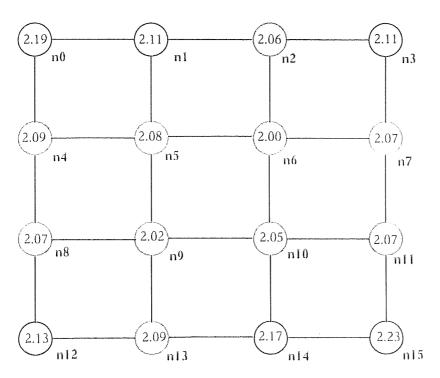
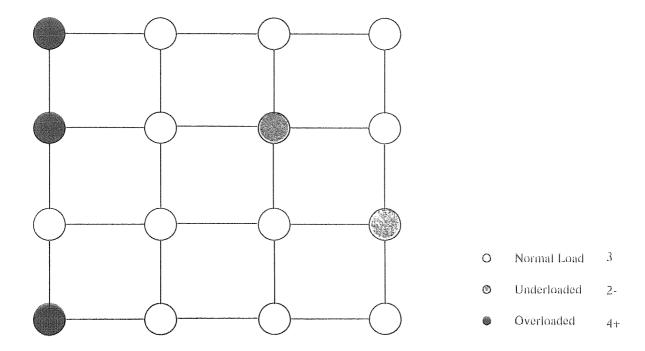


Figure A.12 Response Time After 3000 seconds (Global Average Policy)

Figure A.12 shows the runtime performance after 3000 seconds of simulation time. Whilst the average response time has increased marginally, the run-time performance of nodes 0, 4, and 15 have improved. However, nodes 6 and 9 continue to yield better than average performance.

In terms of their load state after 3000 seconds of simulated time, the following output was produced:

N:	0	Mg:963.00	L:	1	Rt: 2.19	Tx:	12	Imm: 897	Dth: 2460
N:	1	Mg:1283.00	L:	2	Rt: 2.11	Tx:	12	Imm:1242	Dth: 2440
N:	2	Mg:1274.00	L:	2	Rt: 2.06	Tx:	11	Imm:1291	Dth: 2398
N:	3	Mg:992.00	L:	1	Rt: 2.11	Tx:	11	Imm: 992	Dth: 2382
N:	4	Mg:1287.00	L:	4	Rt: 2.09	Tx:	12	Imm:1217	Dth: 2441
N:	5	Mg:1400.00	L:	3	Rt: 2.02	Tx:	13	Imm:1440	Dih: 2416
N:	6	Mg:1340.00	L:	3	Rt: 2.00	Tx:	13	Imm:1418	Dth: 2379
N:	7	Mg:1314.00	L:	3	Rt: 2.07	Tx:	12	Imm:1287	Dth: 2416
N:	8	Mg:1235.00	L:	2	Rt: 2.07	Tx:	12	Imm:1281	Dth: 2371
N:	9	Mg:1414.00	L:	4	Rt: 2.02	Tx:	13	Imm:1442	Dth: 2399
N:	10	Mg:1422.00	L:	3	Rt: 2.05	Tx:	13	Imm:1455	Dth: 2367
N:	11	Mg:1269.00	L:	1	Rt: 2.07	Tx:	11	Imm:1359	Dth: 2317
N:	12	Mg:1055.00	L:	1	Rt: 2.13	Tx:	11	Imm: 944	Dth: 2443
N:	13	Mg:1262.00	L:	3	Rt: 2.09	Tx:	12	Imm:1310	Dth: 2378
N:	14	Mg:1311.00	L:	2	Rt: 2.17	Tx:	12	Imm:1282	Dth: 2428
N:	15	Mg:1027.00	Ĺ:	3	Ri: 2.23	Tx:	11	Imm: 990	Dih: 2401



Overloaded [P(snd) - P(rcv) >> 50]

Underloaded [P(rev) - P(snd) >>50]

Figure A.13 Workload After 3000 seconds (Global Average Policy)

The output indicate that nodes 4, and 10 are overloaded and nodes 0, 1, 2, 3, 8, 11, 12, and 14 underloaded. However, the ratio of the processes received and transmitted identifies nodes 6 and 11 to be underloaded and nodes 0, 4, and 12, to be overloaded. This is illustrated in Figure A.12.

APPENDIX B

PERFORMANCE STATISTICS FOR ALGORITHMS

B.1 The Queuing Model

The graphs in Figure B.1.1 to B.1.10 is representative of the exponential distribution obtained with a total of 100000 processes passing through the system. The load to the system was varied by using the formula proposed by Lavenberg [Lavenberg83] for *traffic intensity* for manipulating the mean interarrival time:

$$\varrho = \lambda E[S] / m \gamma,$$

where $\lambda = 1/E[T]$,

and the notation is explained thus:

λ arrival rate

γ service rate

E[S] Expected service time

E[T] Mean interarrival time

Q system load (stable if $\rho < 1$)

Thus, lengthening the mean interarrival time (giving a low value for ϱ such as 0.2 for example) would represent light system load, and shortening it (such as $\varrho=0.9$) heavy system load. The arriving processes are considered to form a Poisson stream where the interarrival times are independent and identically distributed and are exponentially distributed. The uniformly distributed random number U was accomplished using the UNIX function *erand48* and *drand48*. Both functions are based on the linear congruential algorithm and 48-bit integer arithmetic. The following inverse transformation was then applied:

$$-\lambda^{-1}\log_{e}U$$
,

to form an exponentially distributed random number with mean $1/\lambda$. By using a Poisson process, the system job stream can be split into k substreams, where each substream is Poisson with mean arrival rate λ/k , and exponentially distributed interarrival times. The charts in figures B.1.1 to B.1.10 illustrates this property for both the M/M/1 and M/M/16 queuing models using a Poisson process.

Therefore, the job stream for each host was achieved by first creating k separate data files (jobs 0 _pg0 to jobs $^{k-1}$ _pg0), and generating and splitting a Poisson arrival process using the function erand48 to select the data file for recording the next arrival time. The UNIX function erand48 was also used to create an independent distribution stream for each host in cases where the location policy of the load balancing algorithm needs to randomly select a remote site.

Whilst the results produced by a single simulation run was generally representative of the performance achievable with and without load balancing a total of 20 simulation runs, using different seeds for the pseudo random number generator functions, were conducted per experiment. This was found to be particularly important at high system loads and large processor mesh where performance patterns can exhibit significant fluctuations from one simulation run to another.

Each experiment would run until the average response time for each host differed by less than two percent. Thus, most simulation runs converged to response times of less than two percent after 3600 seconds of simulation time. However, under heavy load conditions, the differences in the performance of edge nodes compared to the more central nodes are re-enforced but stabilise over a longer period of simulated time, namely 8000 to 12000 seconds. Therefore, a 95% confidence interval criteria was used for terminating these experiments. The results produced at low to moderate system load was also found to be consistent with the criteria applied at heavy system load.

Figures B.1.11 to B.1.12, and B.1.13 to B.1.14 shows the average system workload and response times, over the first 1000 seconds of simulated time, for low and heavy system loads respectively. In all cases where load balancing is inactive, the system load and response times stabilise after 600 seconds of simulated time. It is only the global broadcast algorithm that becomes increasingly unstable under extreme load conditions.



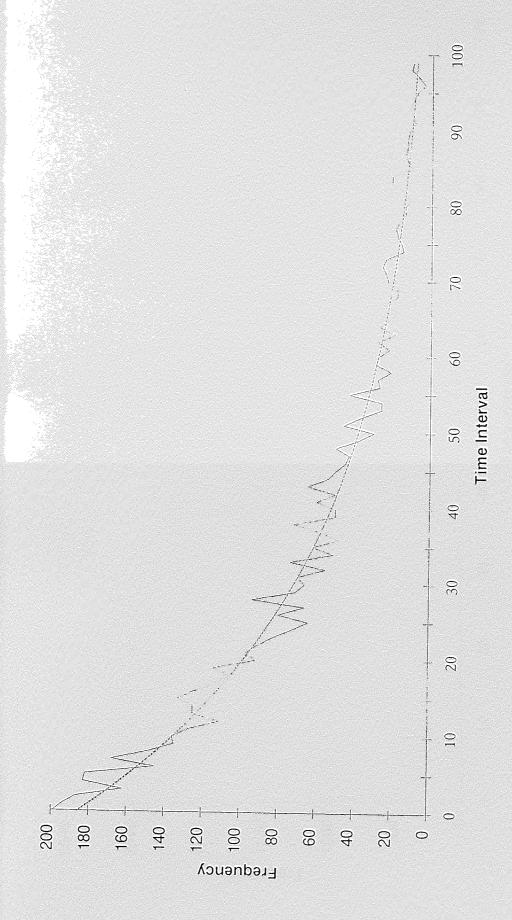


Figure B.1.1 Light Weight Processes (M/M/1): Distribution Curve (Q = 0.2)



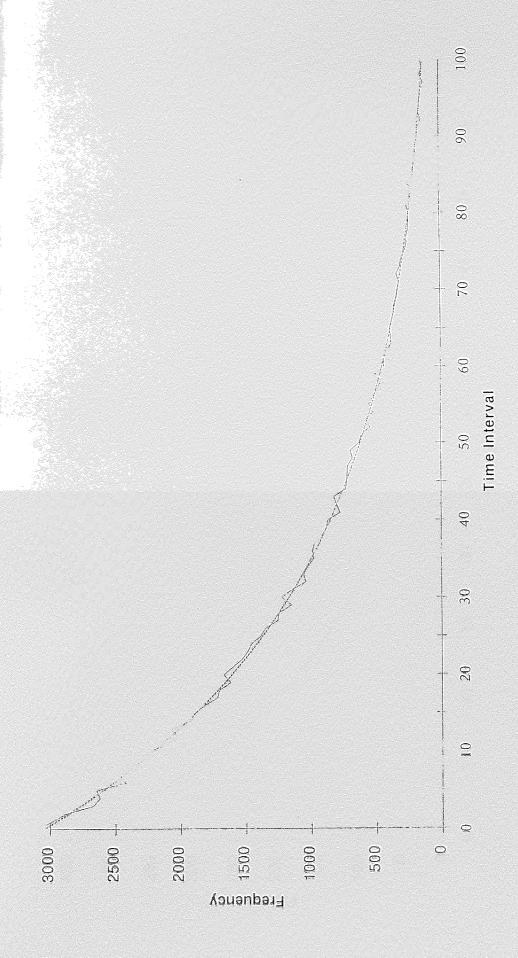


Figure B.1.2 Light Weight Processes (M/M/16) - Distribution Curve (Q = 0.2)

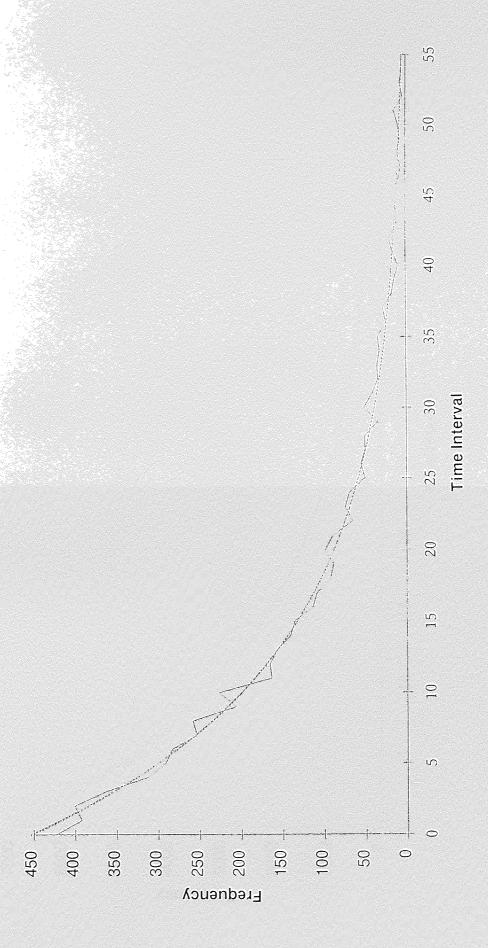


Figure B.1.3 Light Weight Processes (M/M/1): Distribution Curve (Q = 0.5)



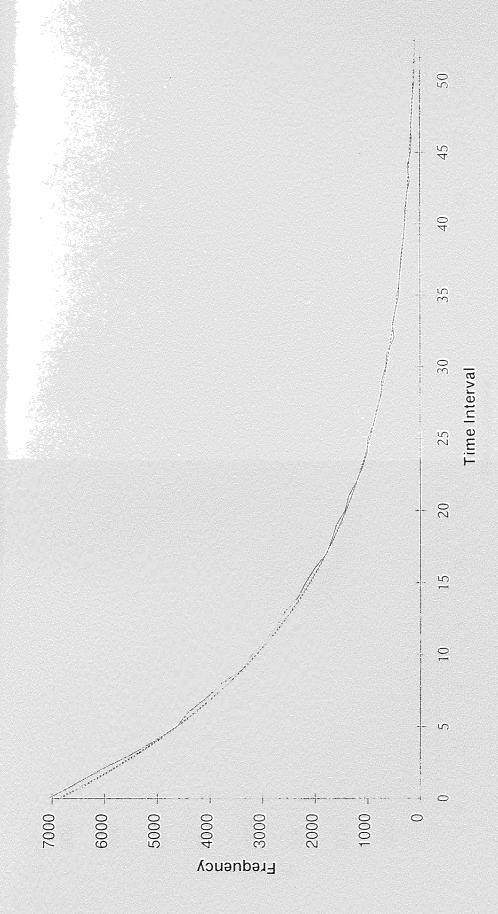


Figure B.1.4 Light Weight Processes (M/M/16): Distribution Curve (Q = 0.5)



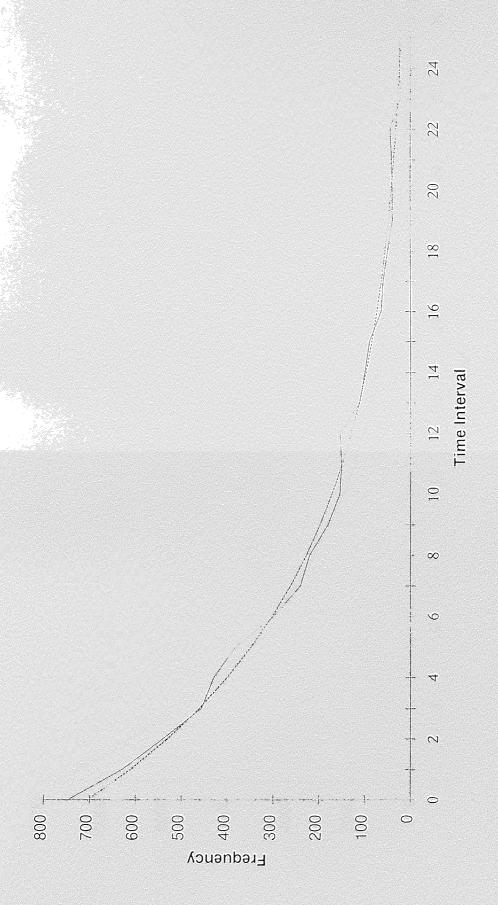


Figure B.1.5 Light Weight Processes (M/M/1) : Distribution Curve (Q = 0.9)



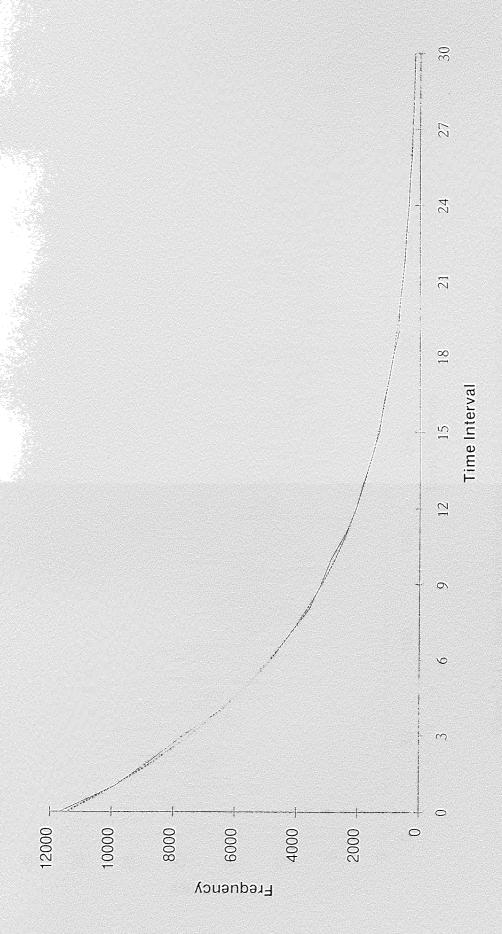


Figure B.1.6 Light Weight Processes (M/M/16): Distribution Curve (Q = 0.9)



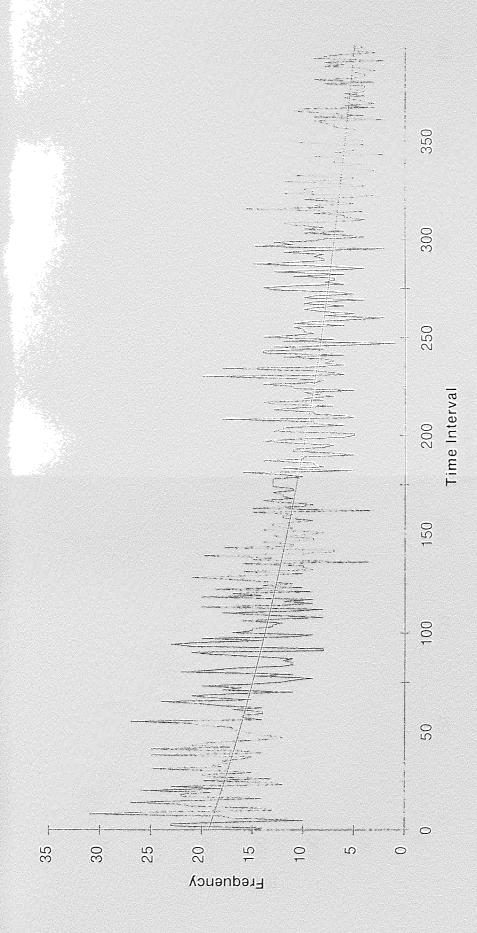


Figure B.1.7 Heavy Weight Processes (M/M/1): Distribution Curve (Q = 0.9)



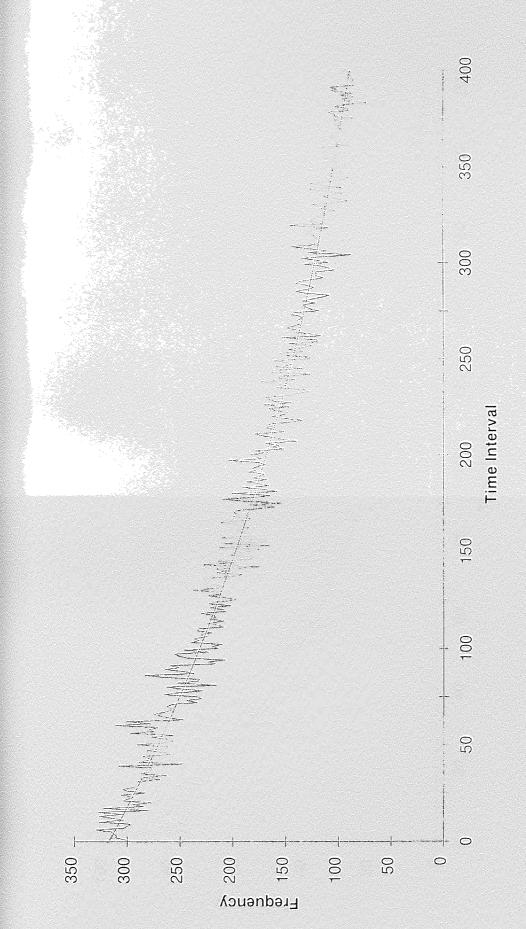


Figure B.1.8 Heavy Weight Processes (M/M/16) - Exponential Distribution Curve (Q = 0.9)



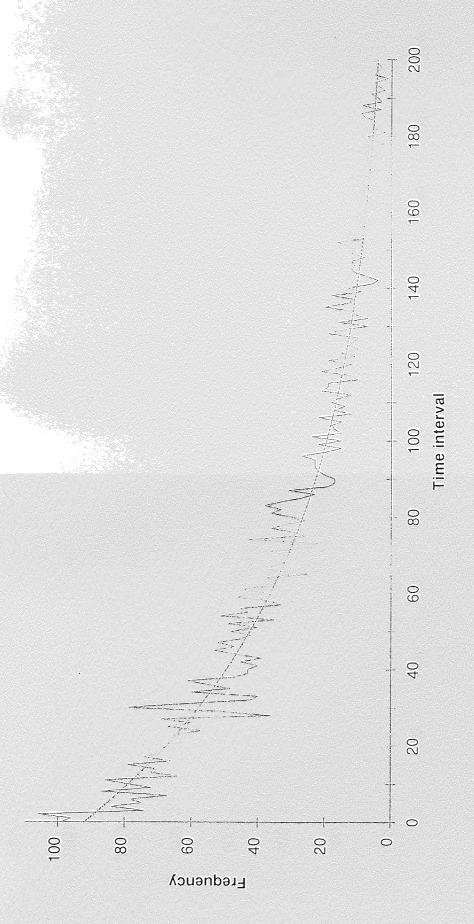
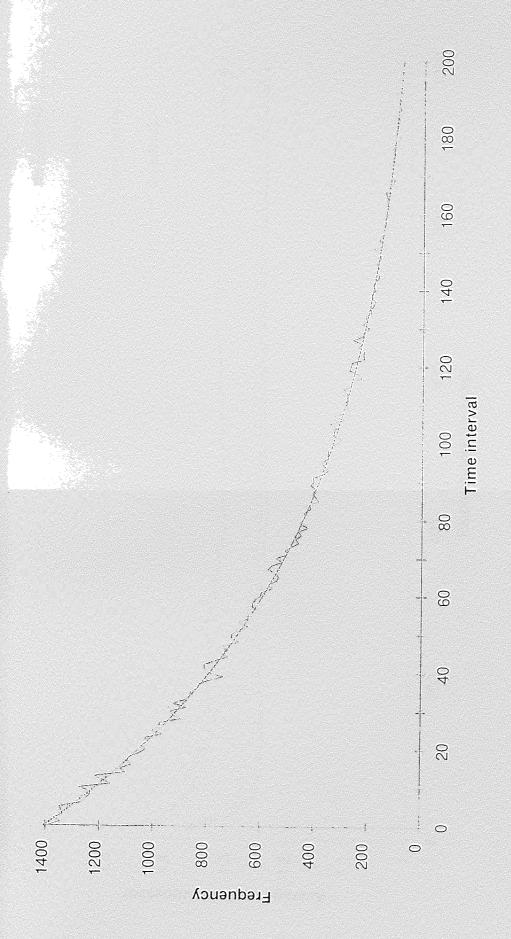


Figure B.1.9 Heavy Weight Processes (M/M/1): Distribution Curve (Q = 0.9)



Figure B.1.10 Heavy Weight Processes (M/M/16) : Distribution Curve (Q = 0.9)



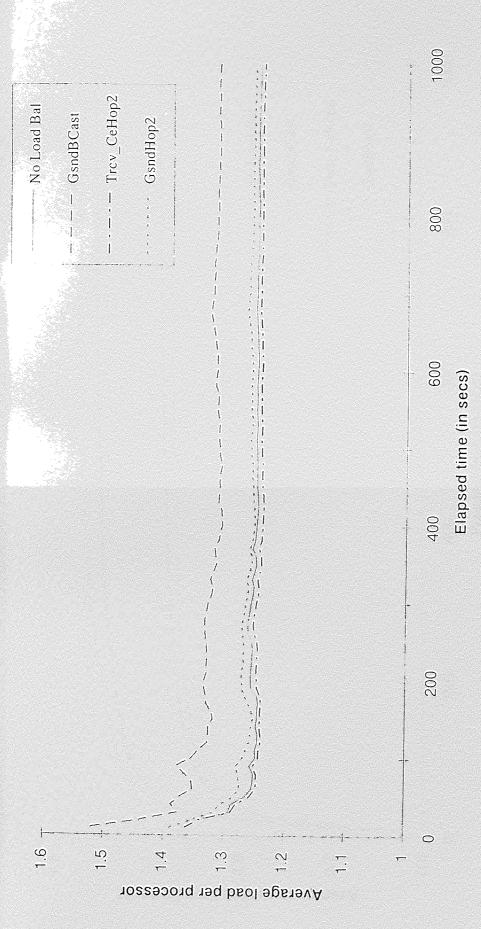


Figure B.1.11 16-Processor Model: Workload Convergence Profile (Q = 0.2)

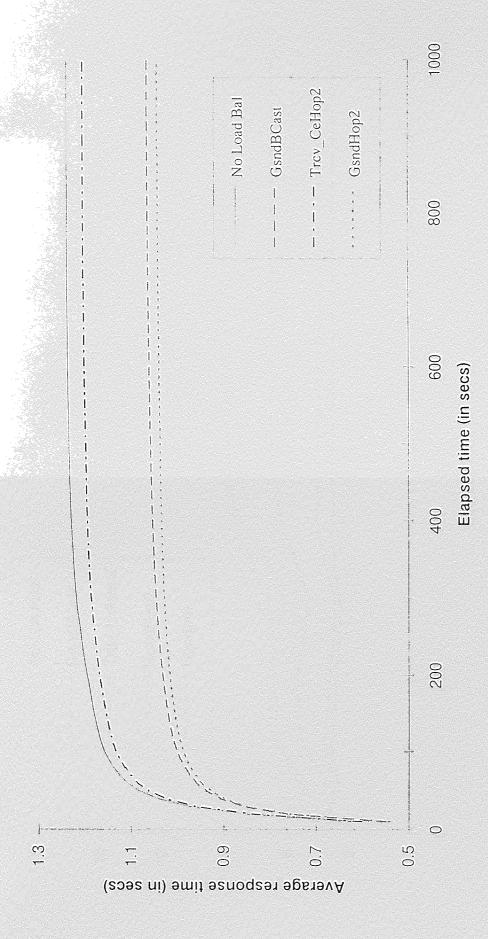


Figure B.1.12 16-Processor Model : Response Time Convergence Profile (Q=0.2)



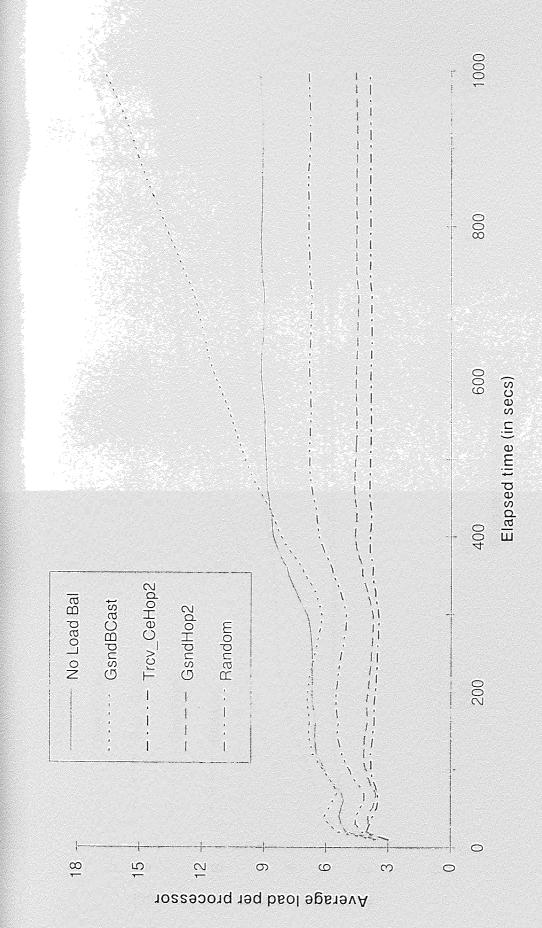


Figure B.1. 13 16-Processor Model : Workload Convergence Profile ($\varrho = 0.9$)

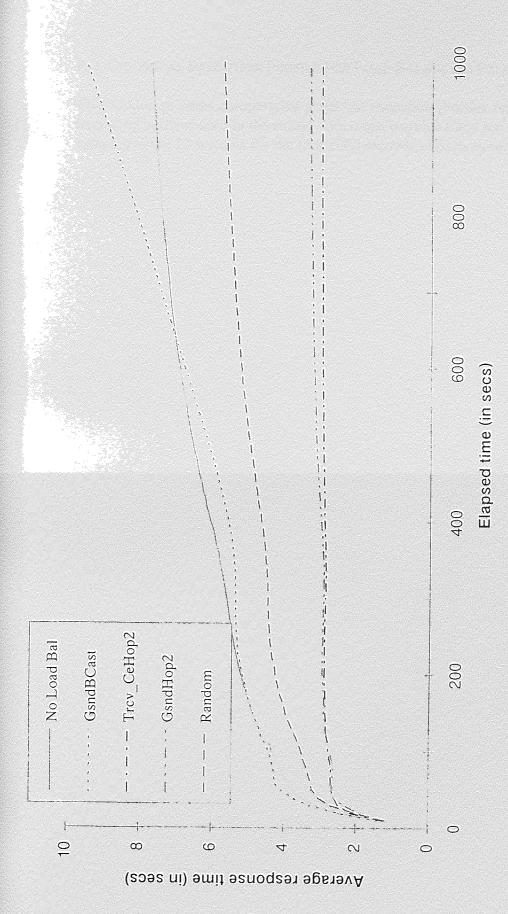


Figure B.1.14 16-Processor Model: Response Time Convergence Profile (Q = 0.9)

B.2 GRAPHICAL EXECUTION PROFILE FOR LOAD BALANCING POLICIES

This section presents comparative graphical execution profiles for a selection of load balancing policies used in the study. The mean workload and run time performance is sampled every 10 seconds for the first 1000 seconds of each simulation run.

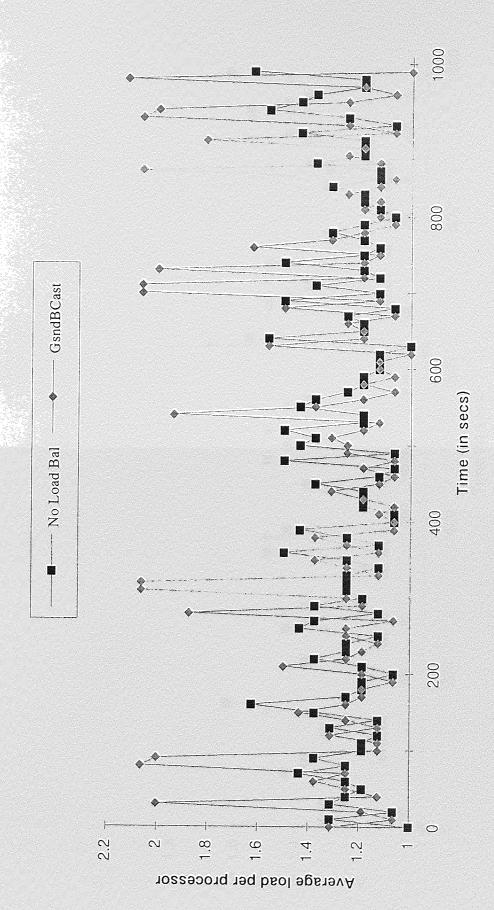


Figure B.2.1 No Load Bal vs Global Broadcast : Lightweight Process Workload Profile ($\varrho = 0.2$)

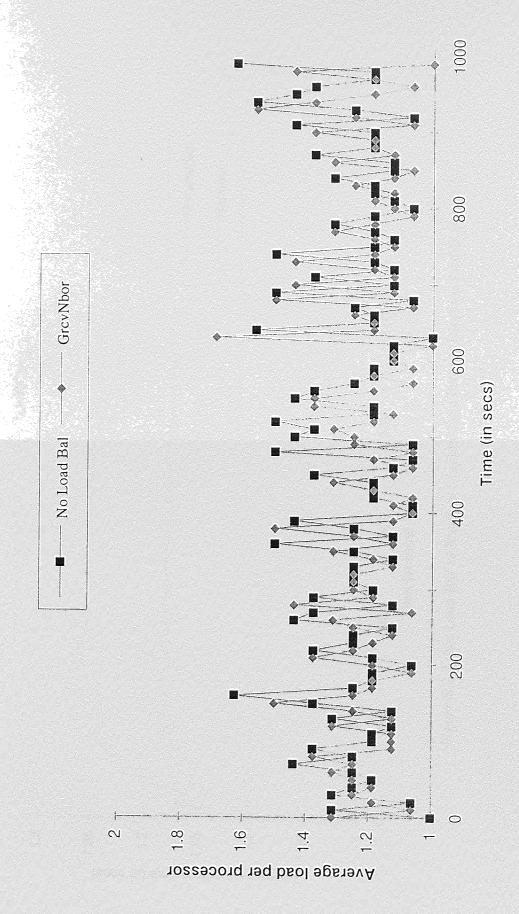


Figure B.2.2 No Load Bal vs Global (rcv) Neighbour : Lightweight Process Workload Profile ($\varrho = 0.2$)



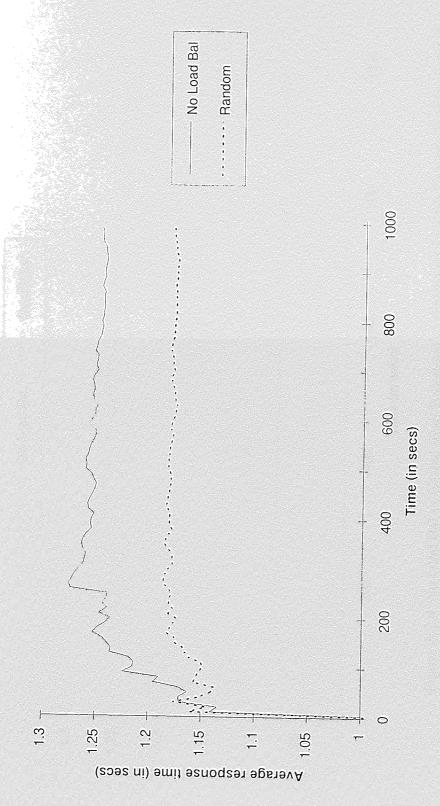


Figure B.2.3 No Load Bal vs Random: Lightweight Process Runtime Profile (Q = 0.2)



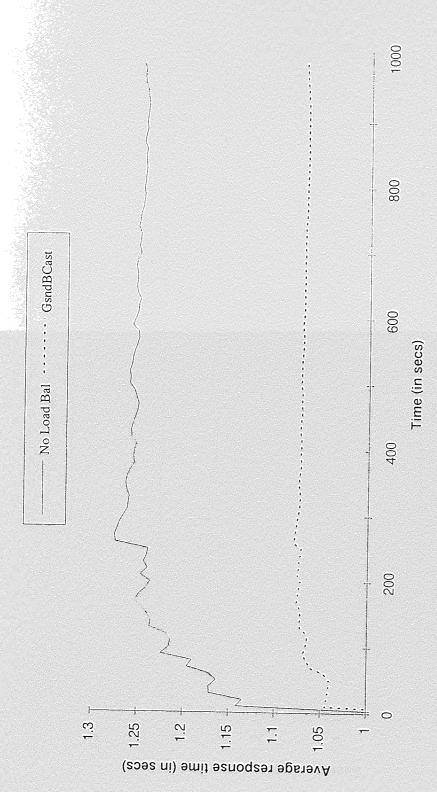


Figure B.2.4 No Load Bal vs Global (snd) Broadcast: Lightweight Process Runtime Profile (Q = 0.2)

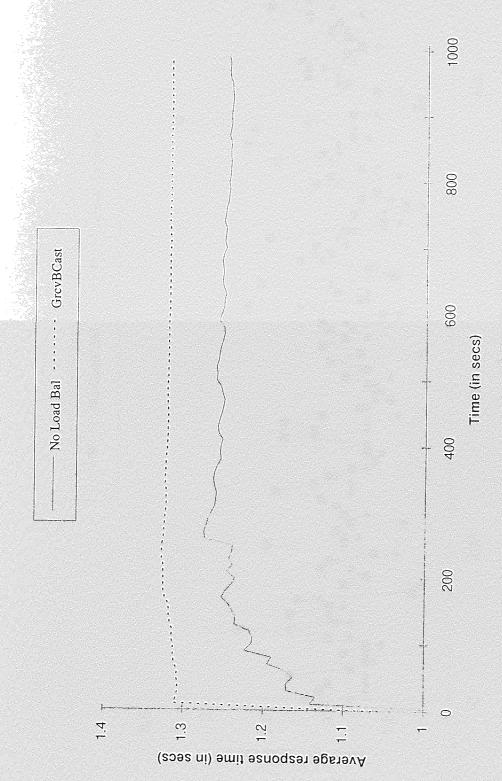


Figure B.2.5 No Load Bal vs Global (rcv) Broadcast: Lightweight Process Runtime Profile (Q = 0.2)

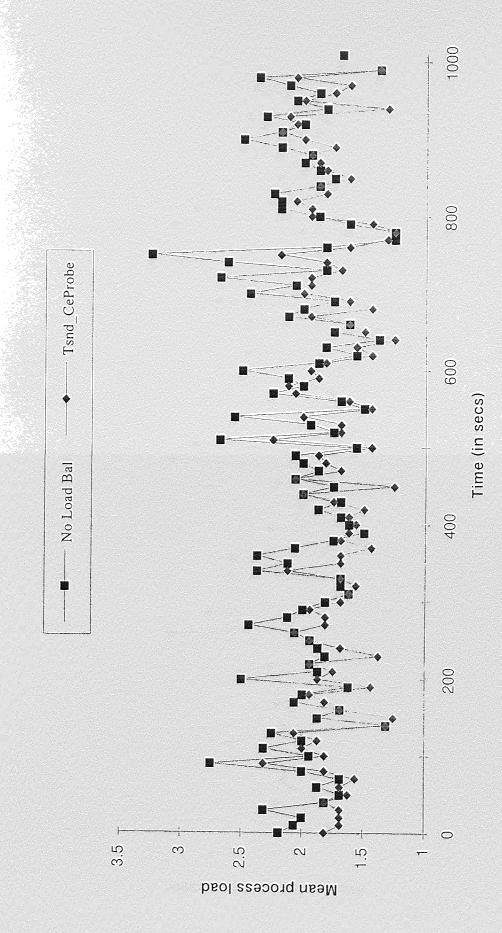


Figure B.2.6 No Load Bal vs Threshold (Cset) : Lightweight Process Workload Profile (Q = 0.5)



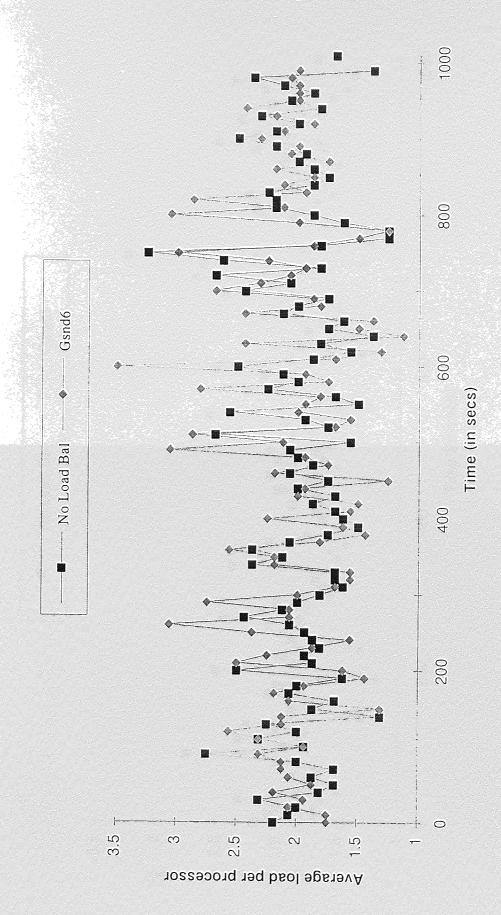


Figure B.2.7 No Load Bal vs Global (snd) Broadcast: Lightweight Process Workload Profile (Q = 0.5)

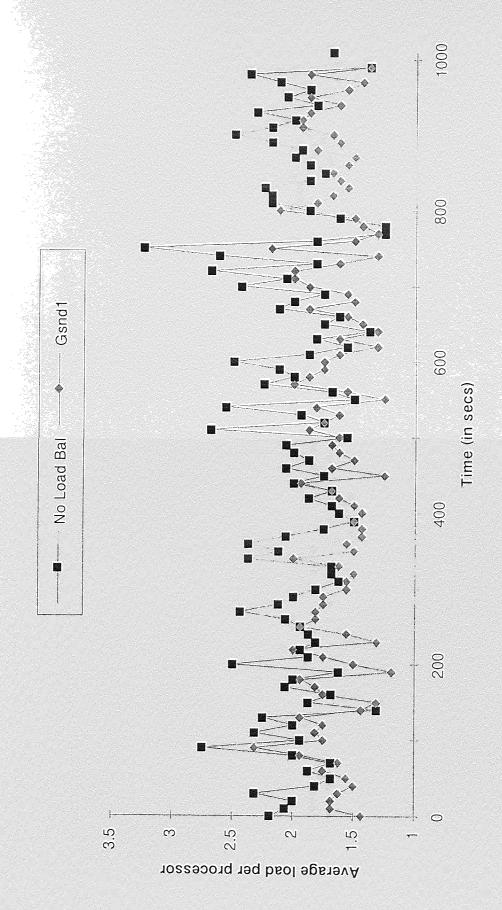


Figure B.2.8 No Load Bal vs Global (snd) Neighbour: Lightweight Process Workload Profile (Q = 0.5)

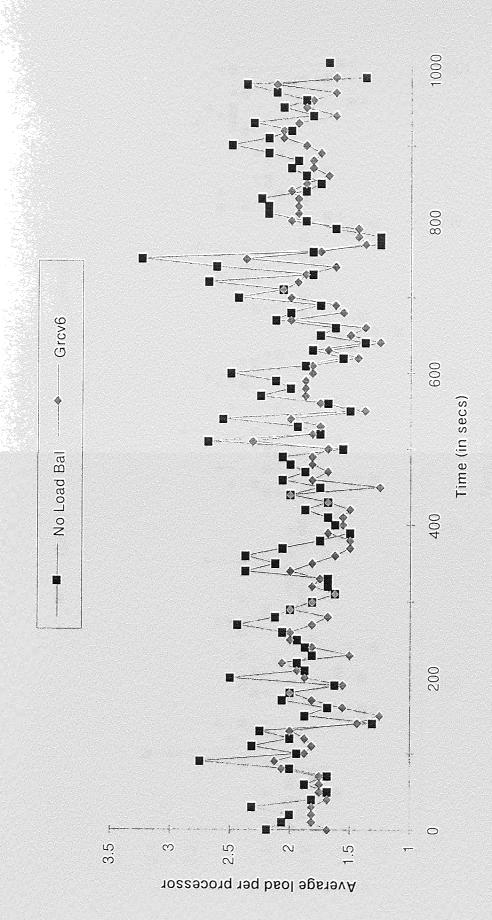


Figure B.2.9 No Load Bal vs Global (rcv) Broadcast: Lightweight Process Workload Profile (Q = 0.5)

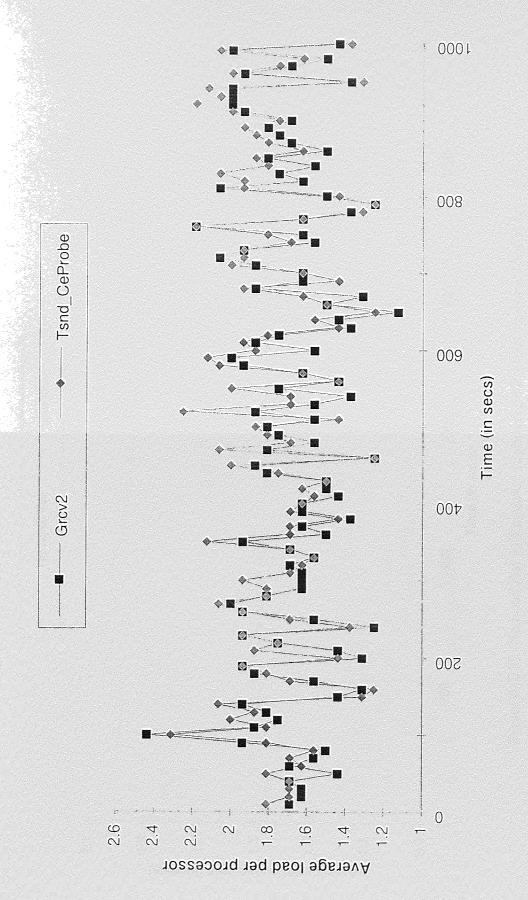


Figure B.2.10 Threshold (Cset) vs Global Receiver (Radius = 2 Hops) : Lightweight Process Workload Profile (Q = 0.5)

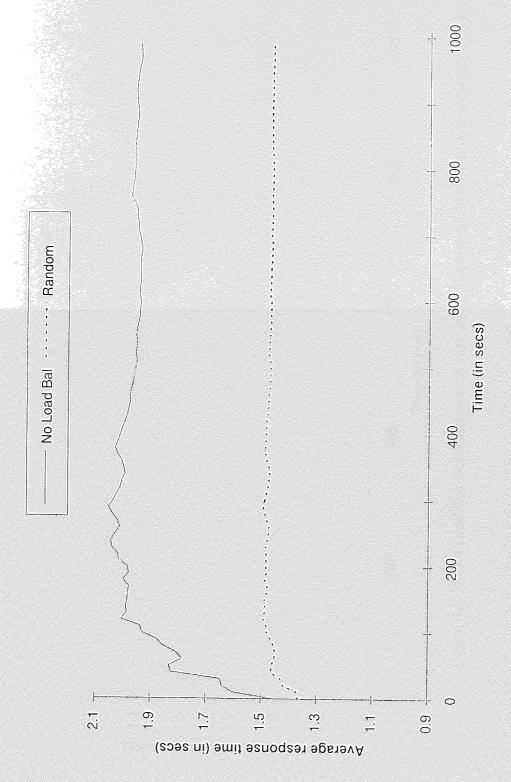


Figure B.2.11 No Load Bal vs Random : Lightweight Process Runtime Profile ($\varrho = 0.5$)

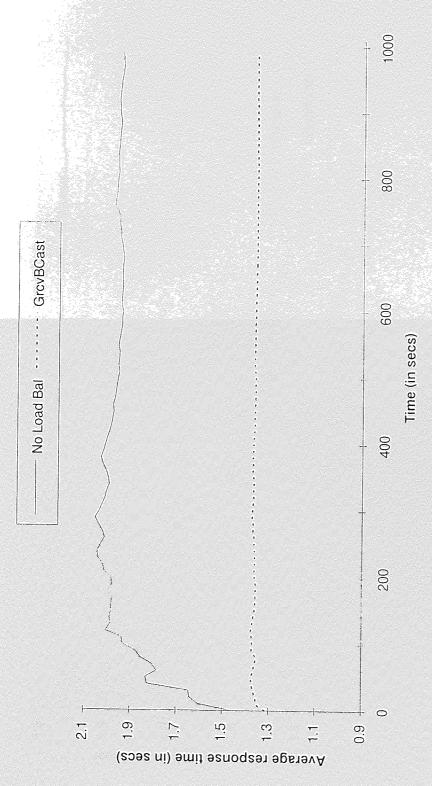
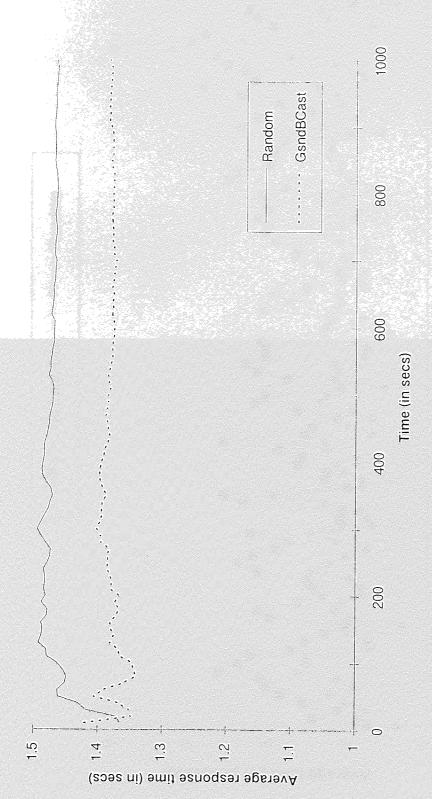


Figure B.2.12 No Load Bal vs Global (rcv) Broadcast : Lightweight Process Runtime Profile (Q = 0.5)





 $Figure\ B.2.13\ Random\ vs\ Global\ (snd)\ Broadcast: Lightweight\ Process\ Runtime\ Profile\ (Q=0.5)$

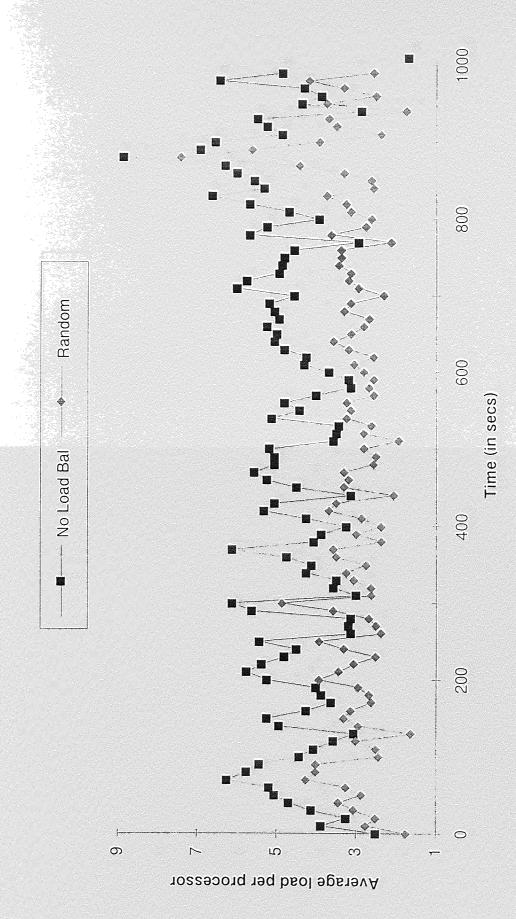


Figure B.2.14 No Load Bal vs Random : Lightweight Process Workload Profile (Q = 0.8)

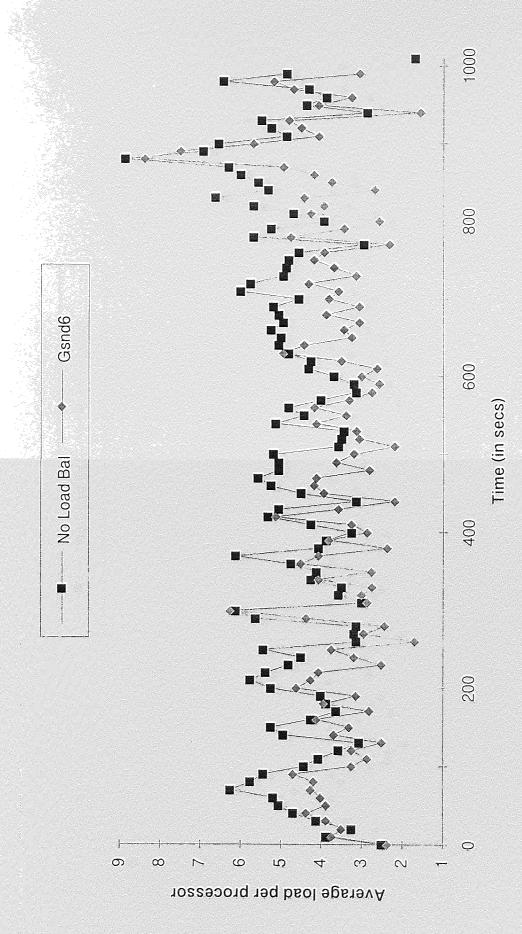


Figure B.2.15 No Load Bal vs Global (snd) Broadcast: Lightweight Process Workload Profile (Q = 0.8)

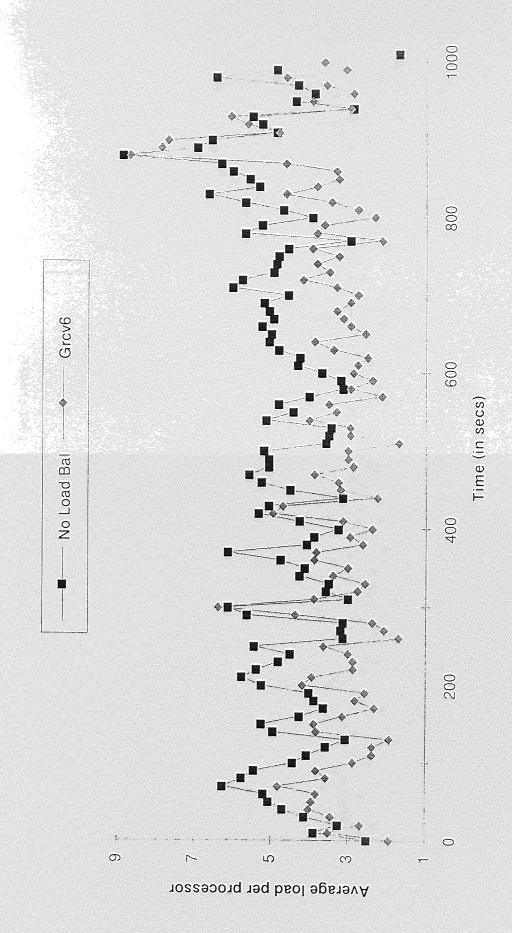


Figure B.2.16 No Load Bal vs Global (rcv) Broadcast: Lightweight Process Workload Profile (Q = 0.8)



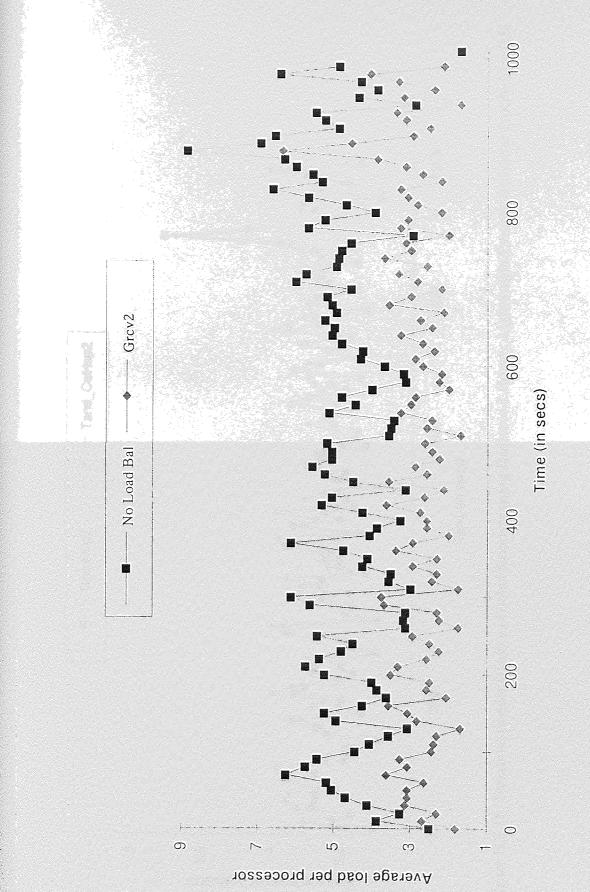


Figure B.2.17 No Load Bal vs Global Receiver (Radius = 2 Hops): Lightweight Process Workload Profile (Q = 0.8)

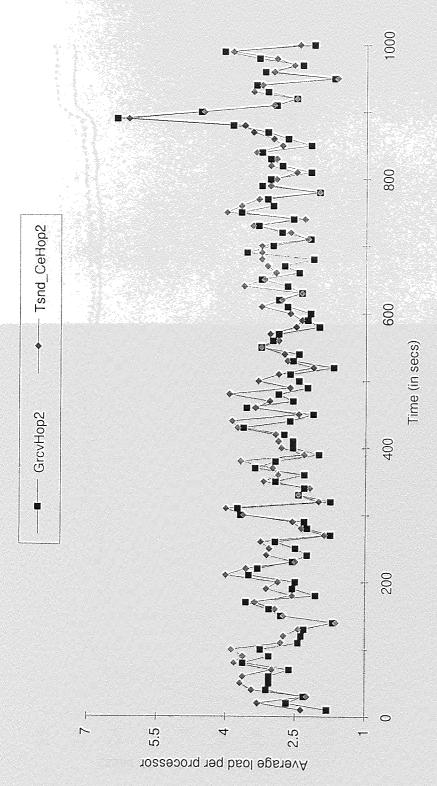


Figure B.2.18 Threshold Sender (Cset) vs Global Receiver using Radius = 2 Hops): Lightweight Process Workload Profile (Q = 0.8)



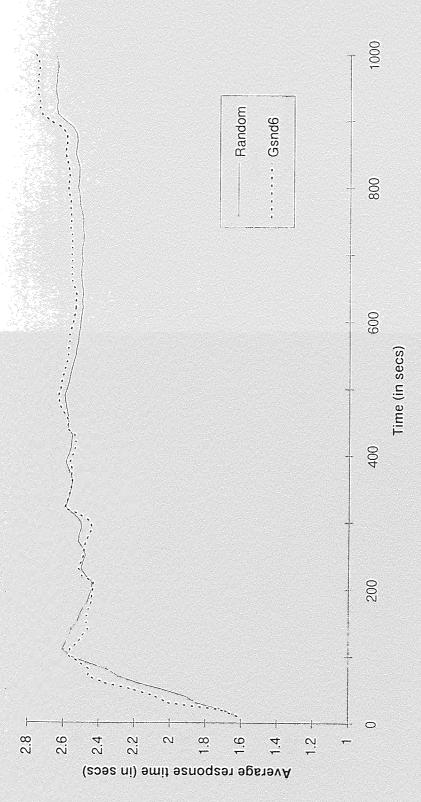


Figure B.2.19 Random vs Global (snd) Broadcast: Lightweight Process Runtime Profile (Q = 0.8)

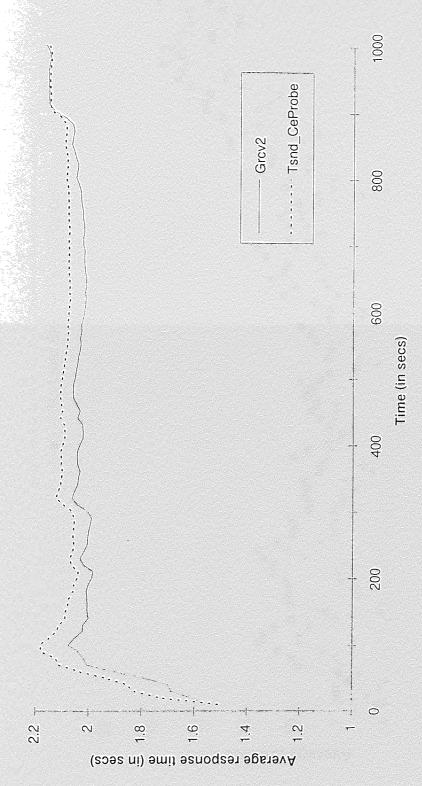


Figure B.2.20 Threshold (Cset) vs Global Receiver (Radius = 2 Hops): Lightweight Process Runtime Profile (Q = 0.8)

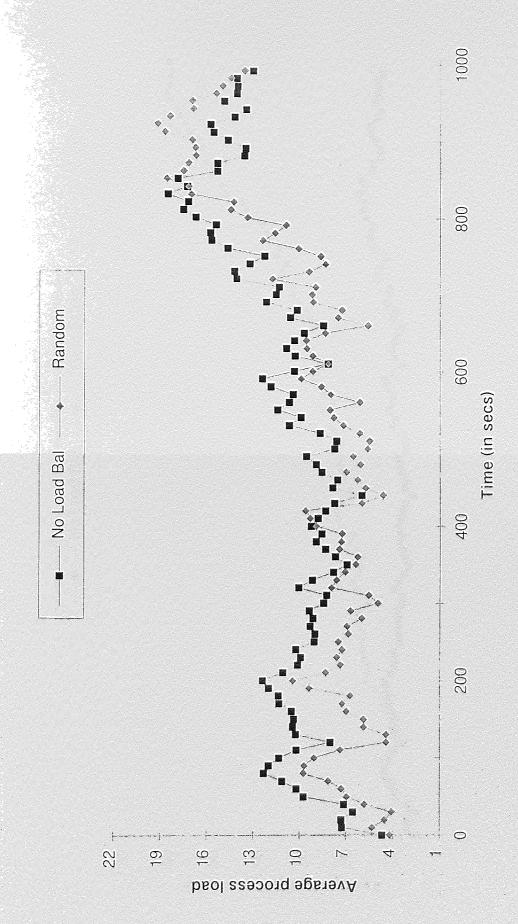


Figure B.2.21 No Load Bal vs Random: Lightweight Process Workload Profile (Q = 0.9)

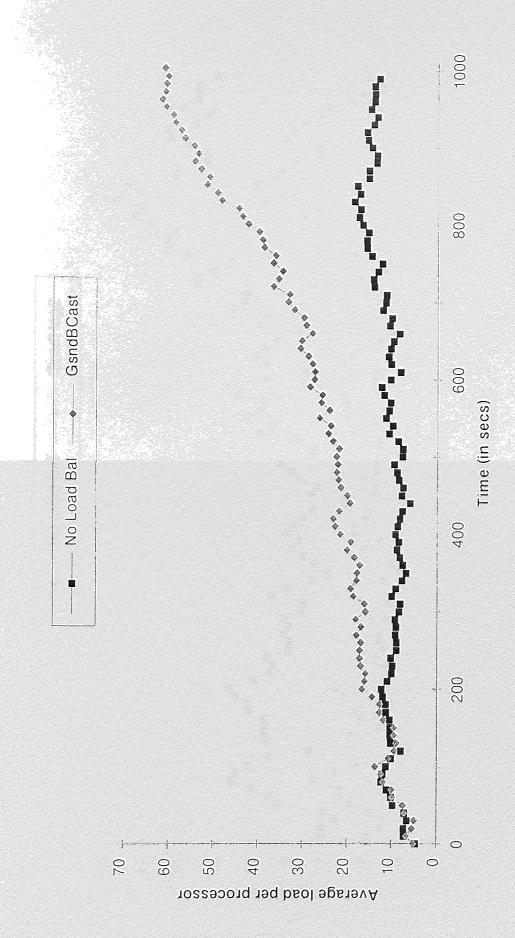


Figure B.2.22 No Load Bal vs Global (snd) Broadcast: Lightweight Process Workload Profile (Q = 0.9)

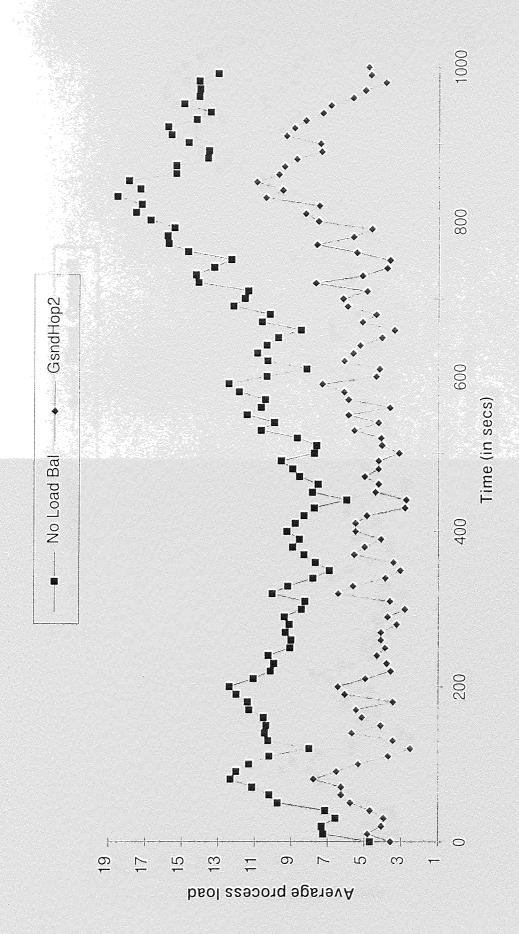


Figure B.2.23 No Load Bal vs Global Sender (Radius = 2 Hops) : Lightweight Process Workload Profile (Q = 0.9)

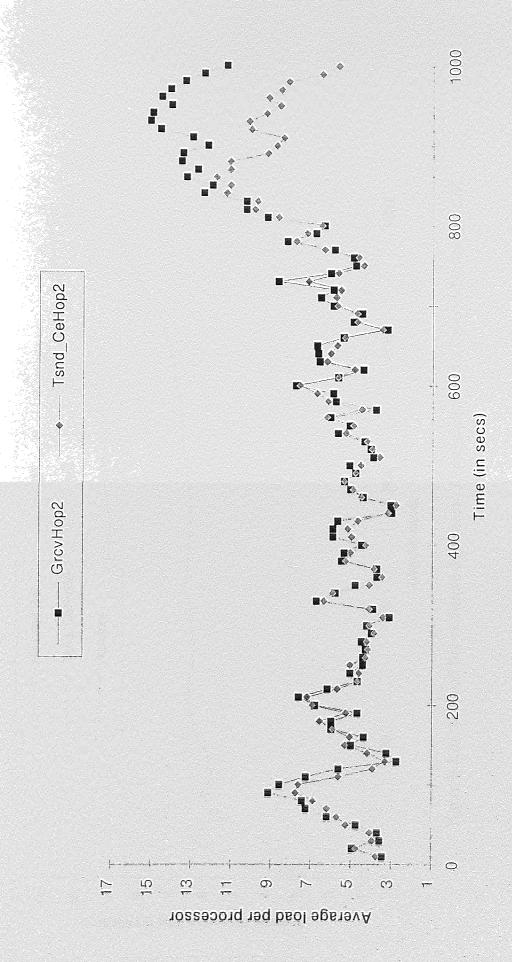


Figure B.2.24 Threshold Sender vs Global Receiver with Radius = 2 Hops: Lightweight Process Workload Profile (Q = 0.9)





Figure B.2.25 No Load Bal vs Simple algorithms : Lightweight Process Runtime Profile (Q = 0.9)



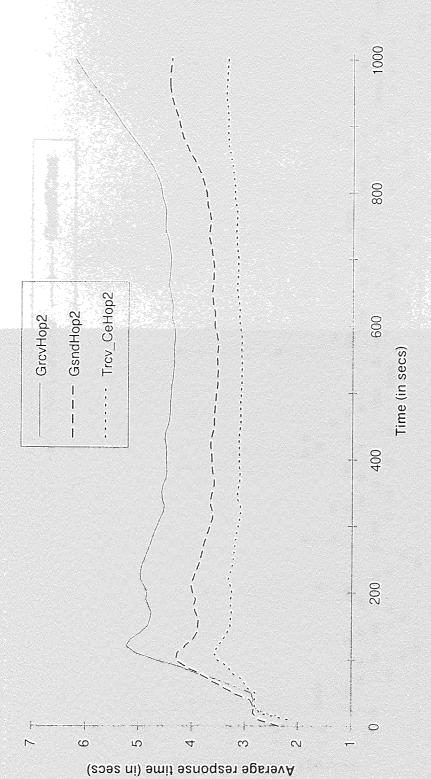


Figure B.2.26 Simple Adaptive vs Complex Global algorithms: Lightweight Process Runtime Profile (Q = 0.9)



200

1

1000

800

GsndBCast

-- No Load Bal

2.4

2.2

<u>~</u> ∞: 1.6

Average load per processor

₽. |4.2



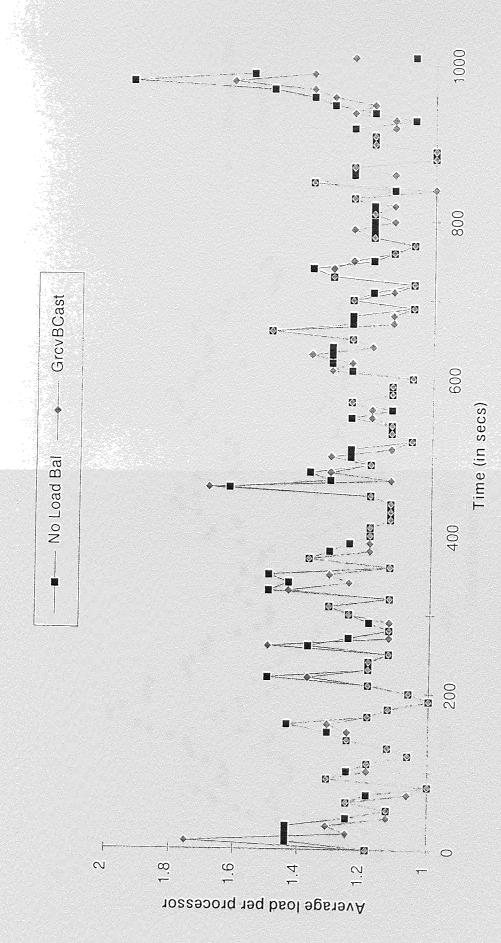


Figure B.2.28 No Load Bal vs Global (rev) Broadcast: Heavyweight Process Workload Profile (Q = 0.2)

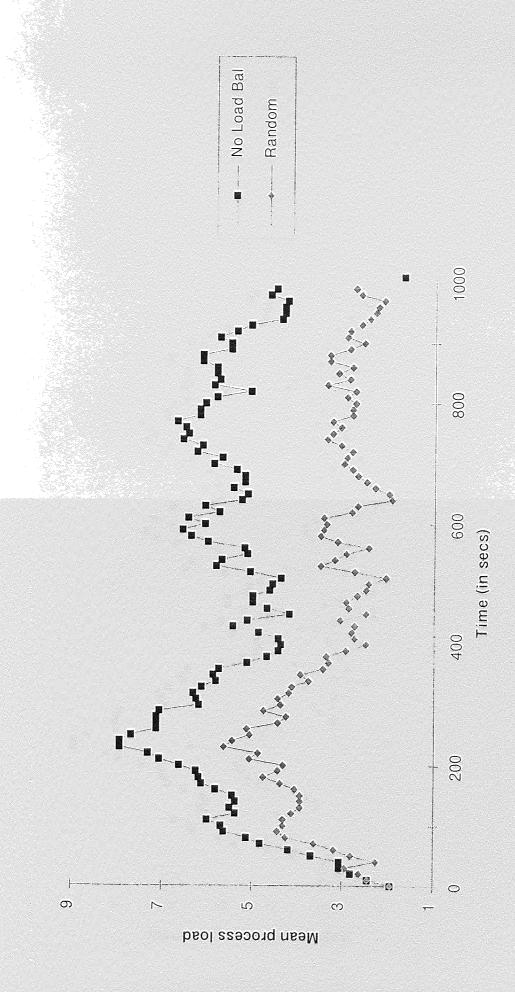


Figure B.2.29 No Load Bal vs Random: Heavyweight Process Workload Profile (Q = 0.8)

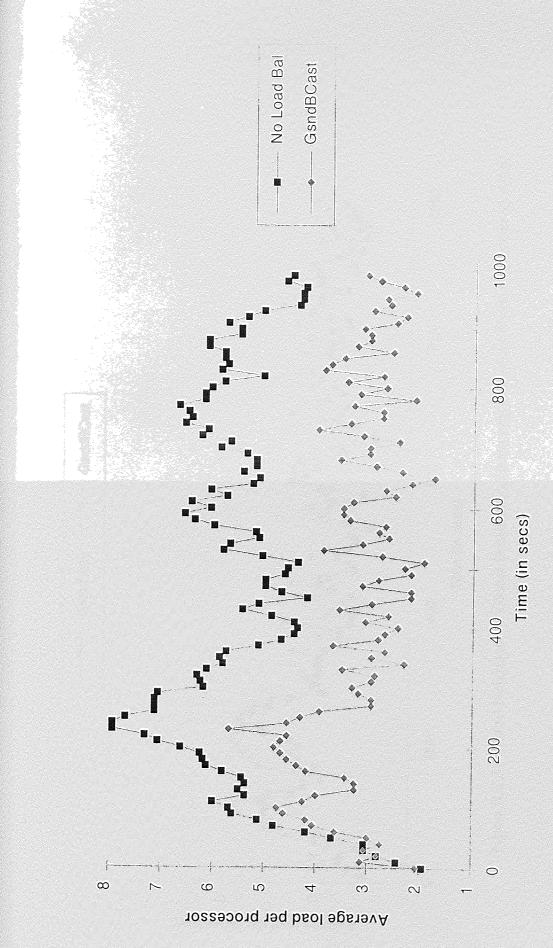


Figure B.2.30 No Load Bal vs Global (snd) Broadcast: Heavyweight Process Workload Profile (Q = 0.8)



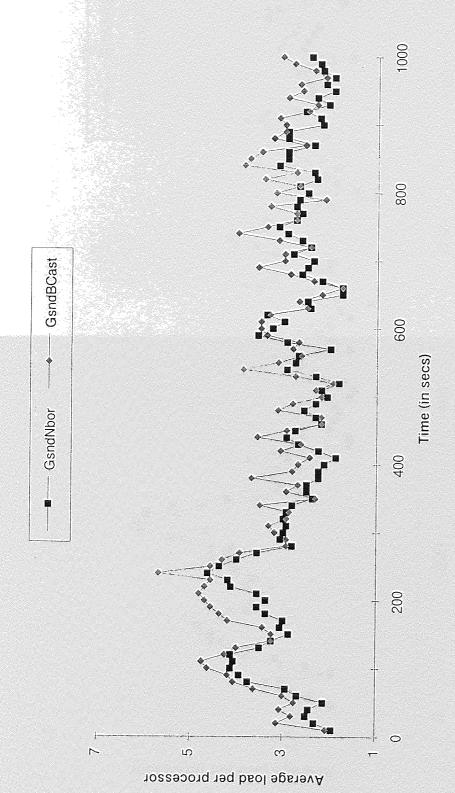


Figure B.2.31 Global Neighbour vs Global Broadcast: Heavyweight Process Workload Profile (Q = 0.8)

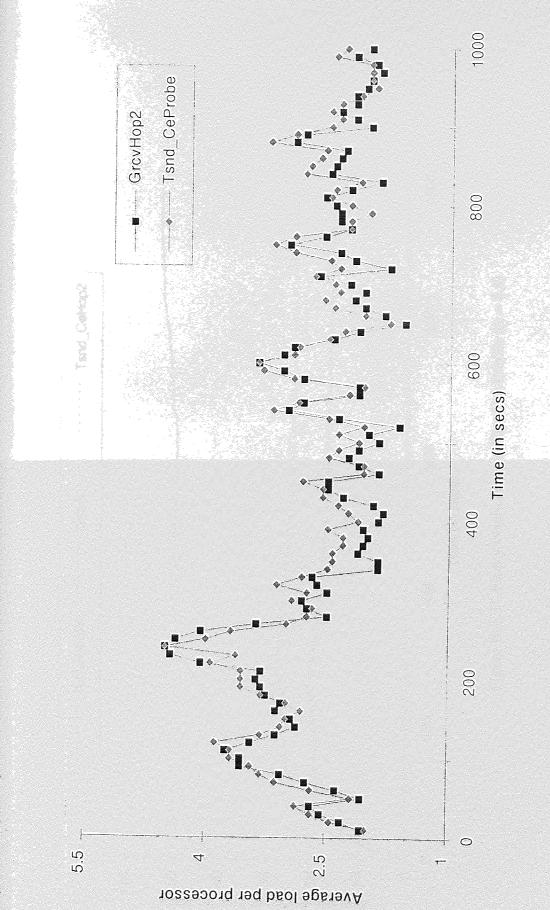
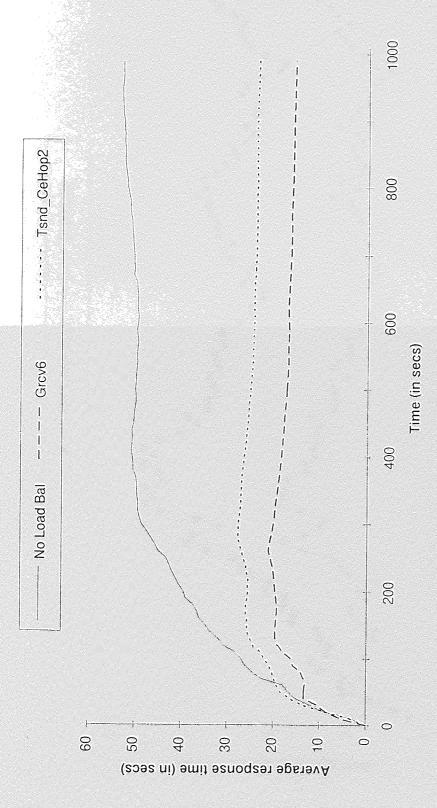


Figure B.2.32 Threshold (Cset) vs Global Sender (Hop = 2): Heavyweight Process Workload Profile (Q = 0.8)



Figure B.2.33 Heavyweight Process Runtime Profile (Q = 0.8)



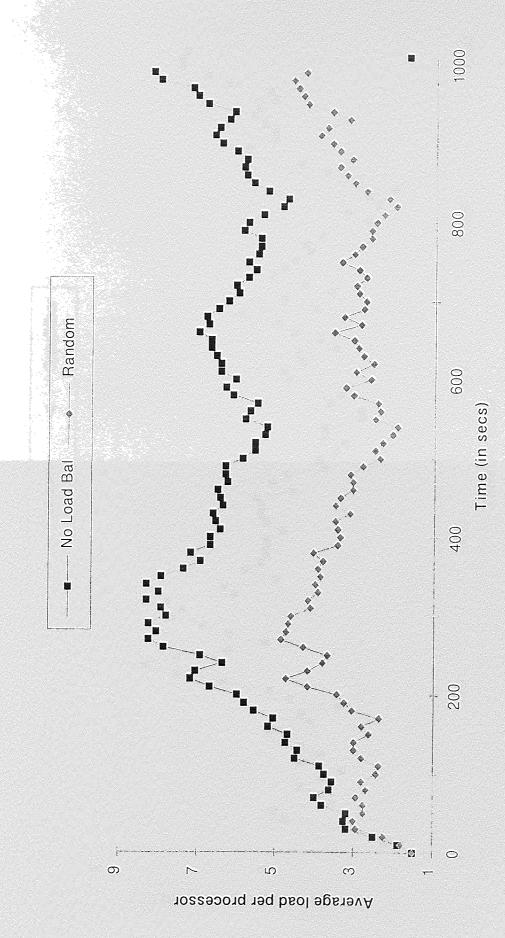


Figure B.2.34 No Load Bal vs Random : Heavyweight Process Workload Profile ($\varrho = 0.9$)

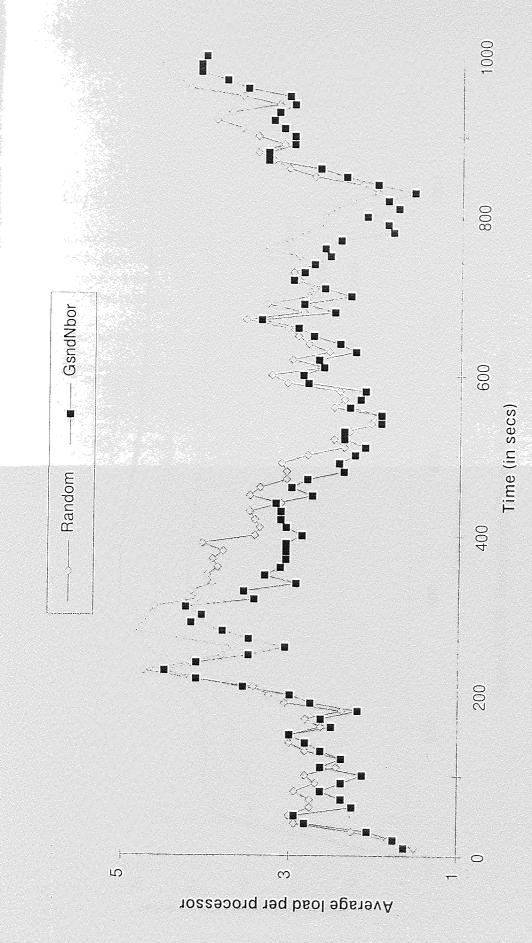
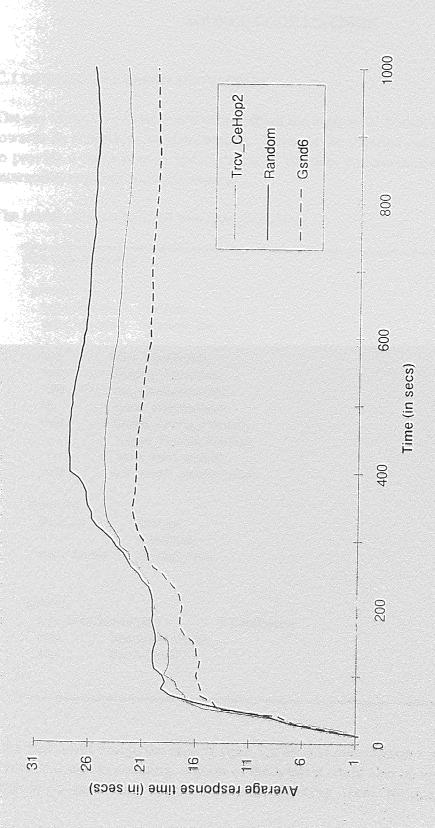


Figure B.2.35 Random vs Global (snd) Neighbour: Heavyweight Process Workload Profile (Q = 0.9)



Figure B.2.36 Heavyweight Process Runtime Profile ($\varrho = 0.9$)



APPENDIX C

SOURCE CODE LISTINGS

C.1 DETAILED C++ DESIGN ABSTRACTION

The use of the C++ class constructs enabled the high-level design to be further refined towards an acceptable implementation. Thus, in the detailed design it was necessary to provide a more precise specification of the data structures, methods, and type of interaction required between objects.

The initial object identified was type_Simulator:

```
class type Simulator
  private:
    type System
                      *system;
    type_Monitor
                     *control_pe;
    char
                *Get_Time(void);
    void
                 Display Menu(int);
    void
                 run_simulation(void);
    int
                 menu selection(int);
    void
                 Create Config File(int),
                 Create_Params_File(int),
                 write intro script();
    void
                 init_model_params (void);
  protected:
   static int
                     total pe,
                     n reached;
   static double
                     G_stop_time;
 public:
   type_Simulator(void);
                                    // Constructor
   void
                Interactive (void);
   void
                Batch (void);
};
```

In the high-level design the objects type_System and type_Monitor were considered as derived classes. In so doing, objects of these types could also be viewed as simulated components providing access functions allowing them to operate in batch or

interactive mode. However, this abstraction was considered invalid as these components may have little in common with type_Simulator when it is used as the base class. Further, an application would need to be developed (of type type_Simulator) to effect simulation using type_System and type_Monitor objects. Therefore, separate classes were adopted for the simulator, the system and the monitor objects. A simulator object can now define system and monitor objects such that a greater degree of control can be exercised over the size of the simulated system and the integrity of any results produced. This is ensured by using private definitions. Thus, if the model was to be represented as UNIX processes, a family group would be used where the simulator object, as parent, creates child processes "System" and "Monitor" and thereafter can order the dialogue with such components using appropriate sequences of IPC calls.

The following application defines a single simulator named *Model* that the application user can run in batch or interactive mode. The application cannot therefore, access directly the time, configuration, and event files of the model.

The following section of code shows the implementation of the *Interactive* method of the simulator object. In particular, a menu-driven interface is used which allows the user to specify the characteristics of the system being simulated. This includes the total number of processors, the network topology, and the load balancing protocol to be used.

Also shown is the implementation of the private function *run_simulation* which activates one or more simulation runs in sequence. Of particular note is the call made to the public functions of the **System** object to *Configure, Boot,* and *Run* the system.

```
void type_Simulator::Interactive(void)
  { int
                menu id,
                new_menu = 1,
                option;
   BOOLEAN
                     exit_loop;
   monitor file = stdout;
   do
   {/* select option and perform function
    menu_id = new_menu;
    option = menu_selection(menu id);
    switch (menu id)
    {/* perform selected function
      case 1: /*
           switch (option)
           {/* perform selected function
             case 1: new_menu = 3;
                  break;
             case 2: new_menu = 2;
                  break;
             case 3: break:
             case 4: run_simulation();
                  break;
             default: break;
          }/* switch */
          break:
      case 2: /*
           switch (option)
           { case 1: break;
             case 4: break;
            default: /* exit from program
                                                 */
                  new menu = 1;
                  break;
          }/* switch */
          break:
     case 3:
          switch (option)
           {/* perform selected function
             case 1: break;
             case 4: break:
            default: /* exit from program
                 new menu = 1;
                 break;
          }/* switch */
          break:
     default:/* exit from program
          break;
    };/* switch */
    exit_loop = (menu_id == 1) && (option == 9);
  } while (!exit loop);
```

```
run simulation
                                                                           */
                                                                           */
                      void type_Simulator::run_simulation (void)
 BOOLEAN
                  positive_response = FALSE;
 int
                    un count = 1,
                   prev_total = 0;
 fprintf (monitor file,
      "\n *** SIMULATION COMMENCED ON %s \n\n", Get_Time());
 trace = fopen("main.trc", "a");
 Create Config File (BATCH RUN):
 Create_Params_File (BATCH RUN);
 configfile = fopen ("config.dat", "r");
 write intro script();
 while (!feof(configfile))
  fscanf (configfile, "%d", &mesh len);
  if (prev total == total pe)
   fscanf (configfile, "%d", &positive_response);
  else prev_total = total_pe;
  system = new type_System(total_pe);
  system->Configure();
  seedfile = fopen("seedfile.dat", "r");
  paramfile = fopen ("pramfile.dat", "r");
  fscanf (paramfile, "%f", &load value);
  while (!feof(paramfile))
   fprintf (monitor_file,
        "\nSIMULATION no.%d STARTED ON %s",
        run count, Get Time());
   init_model_params();
   system->Boot();
   system->Run();
   fprintf (monitor_file,
                         -----\n\n\014");
   fflush (monitor file);
   fflush (trace);
   run count++;
   fscanf (paramfile, "%f", &load_value);
  fclose (paramfile);
  fclose (seedfile);
 system->DisConnect();
 delete (system);
  write_intro_script();
fclose(configfile);
fclose(trace);
```

The class *type_System* contain all functions and methods pertaining to a distributed system, namely, the network topology and the sets of nodes and their characteristics. The class definition is shown below.

```
class type_System
 private:
   void
                     Build_Route Table(void);
   BOOLEAN
                    Balanced_Route (LOCAL_HOST *, int);
   BOOLEAN
                    Find_Links (LINK_ENTRY [],int, int);
                 *Create_ItemQ(int, BOOLEAN, int, LOCAL_HOST *);
   stack entry
   BOOLEAN
                    Stack_Links (LOCAL_HOST *, int);
                    Traverse_Net (LOCAL_HOST *);
   void
   BOOLEAN
                    Push_Ltable (LOCAL HOST *);
   void
               Initialise Netmodel(void);
   void
               Init_Model Params(void);
   void
               Clear Workspace (void);
   void
               Build Network(void),
                System Enable(void);
   float
               Ave Migration(void),
               Ave Transmission(void),
               Mean Difference(void),
               Variance(float),
               Process_Runtime(void);
 protected:
   NODE ENTRY
                      *processor table;
   BOOLEAN
                      Reachable (int, int, int, int);
 public:
   type System(int total pe = 9);
              Dump Tables(void);
   int
             Configure (void);
   void
   void
              Boot (void);
   void
              Run (void);
   void
              DisConnect(void):
};
```

It was at this point that crucial decisions had to be made about a system object and its components. Clearly, a system may be composed of one or more sub-systems. Thus, the individual nodes of a distributed system is a system in its own right, capable of autonomous operation. But, the key difference is their relationship with one another through a separate communication medium. Therefore, it was decided that objects of the type *type_System* class would be considered as "virtual" objects consisting of "real" and "imaginary" components such as the node and the network respectively.

Therefore, the *type_System* class was defined as containing private monitor and processor objects. The constructor for this class is set to a default system size of nine processors, and users of the system object can initiate configuration and reboot of the system as well as "core" (or table) dumps and power down. The following section of code shows the constructor, *Configure*, and *Run* methods for this class

```
type_System::type_System(int total_pe)
{
    processor_table = new NODE_ENTRY[total_pe];
}

void type_System::Configure(void)
{
    Build_Network();
}

void type_System::Run(void)
{
    type_Monitor *control_pe;

control_pe = new type_Monitor;
while (!control_pe->Completed())
{
    System_Enable ();
    control_pe->Update();
}

control_pe->Disable();
delete (control_pe);
}
```

The constructor function creates an array of host objects according to the size of network required. The public access function *Configure* calls the private function *Build_Network* to establish the "logical" links between processor objects.

Note that the monitor object variable (control_pe) is created whenever a system is run and is used in a system object to detected the completion of a simulation run and to update both synchronisation and statistical information.

The class type_Monitor contains all the functions and structures relating to the monitor and control of the system. It is during the design of the monitor class that it became apparent that a monitor object would require access to information deemed to be private to a system such as the group or individual identity of a collection of hosts. In a Network Operating System this is not problematic as all hosts have a public identity known to the user of the system. Alternatively, it could be argued that this information be volunteered by the system object. However, to do so would result in a system object driving the monitor object in a master-slave relationship, thus reducing the effectiveness of the latter. Further, the interface to the system would become more complex in an attempt to make available to the monitor object all the

information it requires. In addition, by making such functions public would result in "open access" to access methods for highly sensitive information.

The alternative would be to make *type_Monitor* a derived class of *type_System*. However, such an abstraction creates other problems as a monitor object would itself be a distributed system inheriting the methods and structures of the *type_System* class. Further, the simplicity and autonomy of the *type_Monitor* class would be lost.

```
class type Monitor
  private:
    void
                    Initialise Nodes(void).
                    Reset_Reached (void);
   float
                    RTime_Convergence(type_System *, int);
    BOOLEAN
                    Stable_System(type_System *);
   void
                    Ave_Compute_or_Display(type_System *, BOOLEAN),
                    Final_Compute(type_System *);
  public:
   type_Monitor(void);
                                  // constructor
   void
                    Update (void);
   BOOLEAN
                      Completed (void);
   void
                    Disable(void):
};
```

It is clear that a monitor object is part of a system and is declared as a private dynamic object to reflect this. Such a declaration should allow the said object privileged access unavailable to other users of the system. However, the C++ implementation does not operate in this manner, but require that each class define who its "friends" are and, as such, have access to private methods and data. Thus, the *type_System* class was refined as follows:

In order that the side effects of global availability be minimised, the above example declared specific methods of the *type_Monitor* class to be friends rather than the complete class. However, this is not ideal as the system object must be identified and passed as a parameteer to the named method of the *type_Monitor* class.

The class *type_LocalHost* was perhaps the most important and difficult class to conceptualise during the design stage, namely because some of its characteristics are inherited from the type of simulation required and the *type_System* class. In a loosely-coupled system, a local host consist of its own operating system and user application processes. Each host will also have its own set of timers and communication "cards".

```
class type LocalHost
    friend class type_System;
   private:
    short int
                       node_id;
    unsigned short
                      xsubi[3];
   LINK_ENTRY
                         *links;
   LOAD_ENTRY
                         cSet[CSET_SIZE];
   int
                  load,
                 threshold,
                 n_tx, n rx,
                 n_migrates,
                 n immigs,
                 n_deaths,
                 n_local_procs,
                 n_active_local_procs,
                 n_virtual procs;
   float
   BOOLEAN
                        high_t_set,
                     low_t_set,
                      pwait set,
                     rplimit;
   ROUTE_ENTRY
                       route_table[TABLE_SIZE];
   double
                       nxt_arr_time,
                     cumul_exist_time;
   type_OPSystem
                      *osVX;
   short int
                     buffer_no;
   char
                    event_fname[16];
   FILE
                     *job file,
                     *trace_file;
  void
                    do_processing (void); /*run */
  void
                    HostNet_Setup(int, int);
  int
                     Ave Call(void);
  BOOLEAN
                       Update_Links (int, STACK_ENTRY *);
  int
                    Dump_MyLinks (int);
  public:
  type_LocalHost(int id = 0);
  void
                    Reset_Pe(void);
  BOOLEAN
                    Suspend Pe(void);
  void
                    Run_Pe(void);
};
```

It is notable that the complete class type_System is declared as a "friend" of the type_LocalHost making the complete components of the latter accessible to the former. In terms of the software of the system, the operating system and user processes were initially considered as independent but related entities. Again, their interrelationship present problems in terms of the proliferation and control of "friends". A further refinement is to consider the user processes as an integral component of a "virtual" software system, namely the operating system. Therefore, the variable osVX represent the operating system interface access to the applications and resources of the local host. The public interface to a local host object consist of three basic functions to reset, suspend and run a local host object. The constructor function for a local host is shown below:

```
*/
     type_LocalHost::TYPE_LOCALHOST
                                                                           */
                                                                           */
type LocalHost::type LocalHost(int id)
{ /* implementation */
    int
         idx;
   /* Construct data node */
    node id
   links
               = new LINK_ENTRY[total_pe];
    for (idx = 0; idx < total pe; idx++)
     links[idx].node id = EMPTY;
     links[idx].hops = total pe;
    for (idx = 0; idx < 4; idx++)
     route_table[idx].node_ptr = NULL;
     route table[idx].node id = EMPTY;
```

This function creates and initialises the links and route tables required for communication purposes. This function also creates (or initialise the operating system object). One can envisage a UNIX shell process being activated at this stage. The implementation of the public functions Run_Pe and $Suspend_Pe$ can then make calls using the public interface of the operating system object. The implementation of these methods follow:

```
void type_LocalHost::Run_Pe(void)
{
  osVX->Msg_Processor();
  if (!Suspend_Pe())
    osVX->Event_Processor();
}
```

```
BOOLEAN type_LocalHost::Suspend_Pe(void)

{
    if (osVX->Time_Check() >= G_stop_time)
    {
        if (!processor_table[node_id].await_sync)
        {
            n_reached++;
            processor_table[node_id].await_sync = TRUE;
        }
        return TRUE;
    }
    else
    return FALSE;
}
```

It is worth noting that a local host will require access to the processor table of the *type_System* class which was achieved by defining "friendships". The following code represents the method to set up the network address tables for the local host.

The operating system is generally regarded as protected software. However, having decided to make user processes a component of the "virtual" operating system object, the issue of kernel and user process characteristics is significant. Whilst such processes may share a common base class, the interrelationship of the kernel processes, from which an operating system is derived, and their interrelationship with user processes created further implementation complexity using the C++ mechanisms. Therefore, the class definition given below identifies the main kernel processes which

include load balancing, local process scheduling, message handling, and time management. A user process object is also a private component accessible using the variable *utask*.

```
class type_OPSystem
 private:
    type_LBProcess
                        *lba:
    type MSGProcess
                        *msg handler;
    class type SCHEDProcess *scheduler;
    type_TIMProcess *clock;
   type_USRProcess *utask;
   float
            Dyn_Calc_Ave_Load(type_TIMProcess *);
   float
            Calc_Ave_Load(type_TIMProcess *);
 public:
  double
              Time_Check(void);
  void
          Msg Processor(void),
           Event Processor(void),
           Restart(void);
           Suspend_Kernel(void);
};
```

In general, user processes and kernel processes have common characteristics. Therefore, a base class *type_AProcess* was defined. Each process has a private process control block, and communicating set information. Likewise, the basic public function available is the return of the process identification number assigned by the operating system. This class definition is given below:

One can then define derived classes type_KNProcess and type_USRProcess using the above base class. However, the primary difference between both process types rests in their priority, the privileges enjoyed, and their behaviour. One had considered

defining processes such as the scheduler and message handler as independent *type_KNProcess* objects, but their interdependency meant that the implementation complexity could be minimised by viewing such processes as private methods belonging to a kernel object. The class template outline is given below:

```
class type_KNProcess:public type AProcess
  private:
// Process Priority store and functions
     void Exit_Process (MSG_FRAME);
     void schedule_process (LOCAL_HOST *);
  protected:
  int
         load.
          threshold,
          n migrates,
          n_immigs,
         n deaths,
          n local procs,
         n active local procs,
         n_virtual procs;
 public:
  BOOLEAN
                  TimeOut_Handler(void);
  void
                  Schedule_Process (LOCAL_HOST *node P);
};
```

Another problem presented by this definition is that the methods of the kernel object must be informed of the associated Host object by means of parameters. Thus, the separate definition of kernel processes as objects of the local host operating system addressed this problem.

The class *type_LBProcess* defined the load balancing methods available. The public interface included service categories for load balancing activity such as sender or receiver-initiated algorithms. The detailed implementation of these algorithms were private to the load balancing process object. The class definition is given below:

```
class type_LBProcess:public type_KNProcess
 private:
    random policy ();
    threshold policy ();
    thresh Bcast ();
    reverse_policy (int);
    Rev avg rcv (int);
    Rev_avg_snd(int);
    global avg snd(int);
    global_avg_rcv(void);
 public:
   receiver_initiated (void);
   sender initiated ();
   Communicating Set (int);
   Virtual_Sender (void);
   Virtual Receiver (void);
};
```

Given the potential number of user processes that can exist in a large system under heavy load conditions, the *type_USRProcess* class was optimised such that every process object that would created in the system did not end up with a set of methods to schedule, create, update and edit itself. Thus, the following class specification consist of a queue of user process control blocks, but possessing the functionality required to emulate a user process:

```
class type_USRProcess:public type AProcess
  friend void type_TIMProcess
      ::Time_Update (type_USRProcess *, int, int);
 private:
  usr PCB
                    *current,
                     *firstib,
                     *lastib;
 public:
  void
             Restart(void):
  usr PCB *
                Add_Process (int, usr_PCB *);
                Find Process (usr PCB *);
  usr PCB *
  usr PCB *
                Remove Process (int);
  usr PCB *
                Create New Process(int);
```

From this definition it is notable that the time_update function of the type_TIMProcess objects needs access to certain components of a user process and, in addition the aggregate class have all the functions that would be required by each user

process to manage itself. The process control block type definition (for usr_PCB) was as follows:

```
struct process_cntrl_block
  int
                upid;
  int
                n probes;
  int
                residency;
  double
               exec time;
  int
                exec_here time;
  int
                residency time;
  double
               exist time;
  int
                orig_mc,
                preferred_mc;
  BOOLEAN
                      schedulable.
               blocked,
               finished,
               mcs_probed[6];
 double
              timeout;
 int
                itype;
 process_cntrl_block
                          *nextP;
```

A similar strategy was also adopted for messages of the system. Thus, as an object each message should be able to manage itself in terms of its construction, send and receive, and any resulting queue positions. Again, the overheads under heavy system load would make such a design impractical. Therefore, the class *type_MSGProcess* was defined with a message queue data structure and the methods required to act upon each element of the structure. These definitions are given below:

```
struct comm_packet
  double
                    time_stamp;
  int
                    sender id,
                    dest_id,
                    distance,
                    total_bytes,
                    type,
                    service no;
  double
                    ldvalue;
  BOOLEAN
                     ack_reply;
  usr_PCB
                    *mess data;
  comm_packet
                     *nextM;
};
```

The class type_SCHEDProcess, given below identifies the a timer object method as "friendly" and, in the public interface, the scheduling discipline to be operated. The class definition is given below:

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Likewise, the class *type_TIMProcess* definition is given below. It is notable that all time-dependent and related functions are included in this template. These include timeout message handlers. The code displays some of the weaknesses of a C++ implementation using "friends" and object parameters for autonomous but related classes.

```
class type_TIMProcess:public type_KNProcess
  friend float type_OPSystem::Dyn_Calc_Ave_Load(type_TIMProcess *);
  friend float type_OPSystem::Calc_Ave_Load(type_TIMProcess *);
  private:
   double
                    sys_real time,
                    next_elapsed second,
                    last_perf_dump_time;
  QTUM_ENTRY
                            quanta[NQUANTA];
  int
                            quanta idx,
                            OSoverhead;
  TMER QUEUE
                            timq;
                    Wrq_Timout_Msg(void),
    void
                    Hi_Timeout_Msg(void),
                    Avgld_Tout_Msg(void),
                    Low_Timeout Msg(void);
    void
                    Select_Next Timeout(void);
    void
                    Service_Routine(void);
    void
                    Quanta Update(int, int);
    void
                    Next_Quantum(void),
                    Add_To_Quantum(int, int);
  public:
    BOOLEAN
                    Cancel_Timer (int, usr PCB *);
    BOOLEAN
                    TimeOut Handler(void);
    void
                    Restart_Clock(int , double, usr_PCB *);
    void
                    Reset(void),
                    Time_Update (type_USRProcess *, int, int);
    double
                    Display_Time(void)
                    {return sys_real_time;}
};
```

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Whilst it is good object-oriented programming practice to avoid using "friends" and optimising the behaviour of objects, attempting the implementation of the model in C++ was very limiting. The proliferation of "friends" and the passing of objects as parameters increased the complexity of the interface between objects, and reduced the potential size of the system that could actually be modelled. Therefore, the high level design and the methods developed were grouped into five ANSI C source files which elliminated "friendships" and provided for a more efficient and meaningful implementation. The ANSI C listing for the model is given in the next section. The files are:

The design exercise was not fruitless as it enabled an early identification of all the methods that would be required in the implementation. Given a more flexible object oriented implementation language that retain the links between high-level abstractions and detailed implementation, the software construction stage for the model would have contained the actual amount of software developed for the experiments conducted.

The remaining portion of this chapter gives a sample of some of the implementation methods prototyped for the classes specified.

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```
void type_TIMProcess::Reset(void)
  sys_real_time = 0.0;
  last perf dump time = 0.0;
  next_elapsed second = 1000000.0;
  quanta idx
                = 0;
  OSoverhead
                 = 0;
  for (int idx2 = 0; idx2 < NQUANTA; idx2++)
     quanta[idx2].actual load = N SYS PROCS;
     quanta[idx2].virtual_load = 0;
     quanta[idx2].OSportion = 0;
     quanta[idx2].used
  timq.minP = NULL;
  timq.itype = 0;
  timq.mintime = 0;
```

```
void type_USRProcess::Restart(void)
{
    current = NULL;
    firstjb = NULL;
    lastjb = NULL;
}

void type_MSGProcess::Restart(void)
{
    queue.tail = NULL;
    queue.head = NULL;
}

void type_LocalHost::Run_Pe(void)
{
    osVX->Msg_Processor();
    if (!Suspend_Pe())
        osVX->Event_Processor();
}
```

```
BOOLEAN type_LocalHost::Suspend_Pe(void)
  if (osVX->Time_Check() >= G_stop_time)
   if (!processor_table[node_id].await_sync)
     n reached++;
     processor_table[node_id].await_sync = TRUE;
   return TRUE;
  else
   return FALSE;
void type_OPSystem::Msg_Processor ()
 BOOLEAN rcv;
 do
  rcv = msg_handler->Receive();
  msg_handler->Process_Message();
  clock->TimeOut_Handler();
  if (!rcv)
    break;
 } while (1);
```

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C.2 SHELL SCRIPTS FOR DATA MANAGEMENT

This section contain shell scripts that were written to convert the trace output of the simulation model into a format suitable for performing statistical analysis. Thus, the following output is typical for a collection of simulation runs:

*** SIMULATION COMMENCED ON Tue Jun 1 17:26:23 1993

Load Balancing algorithm invoked every 20 - 50 milliseconds!

SIMULATION no.1 STARTED ON Tue Jun 1 17:26:23 1993

No. of Processors = 9 Load value = 0.80

Total No. of jobs = 1000000 Ld Balancing Alg. = Global Avg Nbor Policy

Threshold Level = 2.00 Distance = 4 Convergence to < 2.00%

TIME VAR. LOAD DIFF. RTime %Conv MIG. Τx %JOB 1800.00 0.7850 3.9543 2.4406 3.3098 0.8988 36.0651 1.29 7.48 0.7792 3.9503 2.4364 3600.00 3.1571 2.85 0.8968 35.8049 2.58 *** TERMINATING CONDITION: CONVERGENCE @2.00%***

SIDIABSE BUILDHONE

SIMULATION COMPLETED ON Tue Jun 1 17:32:24 1993

SIMULATION no.2 STARTED ON Tue Jun 1 17:32:24 1993

No. of Processors = 9 Load value = 0.80

Total No. of jobs = 1000000 Ld Balancing Alg. = Global Avg Nbor Policy

Threshold Level = 2.00 Distance = 1 Convergence to < 2.00%

VAR.	LOAD	DIFF.	RTime	%Conv	MIG.	Tx	%JOB
1.2337	2.8964	2.9672	2.2577	3.20	0.6477	8.1036	1.29
1.2880	2.9055	3.0283	2.2175	0.88	0.6530	8.1900	2.58
1.2867	2.9041	3.0266	2.2175	0.94	0.6529	8.1880	2.59
	1.2337 1.2880	1.2337 2.8964 1.2880 2.9055	1.2337 2.8964 2.9672 1.2880 2.9055 3.0283		1.2337 2.8964 2.9672 2.2577 3.20 1.2880 2.9055 3.0283 2.2175 0.88	1.2337 2.8964 2.9672 2.2577 3.20 0.6477 1.2880 2.9055 3.0283 2.2175 0.88 0.6530	1.2337 2.8964 2.9672 2.2577 3.20 0.6477 8.1036 1.2880 2.9055 3.0283 2.2175 0.88 0.6530 8.1900

*** TERMINATING CONDITION: CONVERGENCE @2.00%***

SIMULATION COMPLETED ON Tue Jun 1 17:36:01 1993

SIMULATION no.3 STARTED ON Tue Jun 1 17:36:01 1993

No. of Processors = 9 Load value = 0.80

Total No. of jobs = 1000000 Ld Balancing Alg. = Global Avg Nbor Policy

DIFF.

Threshold Level = 2.00 Distance = 2 Convergence to < 2.00%

LOAD

VAR.

TIME

RTime

%Conv

MIG.

Tx

%JOB

What is required is to extract the convergence results for each load balancing policy, computing the duration of the simulation run and the seeding used for the random number generator. The following script command would achieve this objective:

xcel run_file [No_Alg]

The parameters required are the name of the run file and the number of load balancing policies being used. Thus, the above output would be transformed as follows:

3600.00	0.7792	3.9503	2.4364	3.1571	2.85	0.8968	35.8049	2.58			
Global Average Neighbour Policy/Distance 4											
3610.00	1.2867	2.9041	3.0266	2.2175	0.94	0.6529	8.1880	2.59			

Global Average Neighbour Policy/Distance 1											

Global Average Neighbour Policy/Distance 2											

The resulting file can then be ported into an spreadsheet package (such as EXCEL) where statistical analysis can then take place. The listings below are of all the scripts required to produce the above output.

```
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```

RTIME_AWK

```
BEGIN {
    value1 = ""
    value2 = ""
     value3 = ""
     value4 = ""
     value5 = ""
     son = ""
 if (\$7 == "Ld")
  value3 = $11 " " $12$13$14
 if ($5 == "Probe")
  value5 = $8
 if ($5 == "Distance")
  value5 = $7
 if ($2 == "TERMINATING")
  son = " "value2 " "value3 value5 "\t"
 else
 if ($1 == "Final")
  value4 = substr(\$3,1, index(\$3,",") -1) "\t" \$6
  print son value4
  value2 = value1
  value1 = $0
}
```

```
BEGIN {
     value2 = $0
     AccTime = 0
     count = 0
     name = ""
     prevname = ""
 if (idx \% algs == 0)
  print value2
  count++
  AccTime += $5
  name = prevname
 idx++
 prevname = $10$11
 value2 = \$1"\t"\$2"\t"\$3"\t"\$4"\t"\$5"\t"\$7"\t"\$8
 value2 = value2 \ "\t"\$14"\t"\$6"\t"\$9"\t"\$12"\t"\$13"\t"\$15"\t"\$16"\t"\$17
END {
if (idx % algs == 0)
 print value2
 print "\n AvgRT = "AccTime" / "count " = " AccTime/count"\t"name"\n"
```

```
#include "stdio.h"
 #include <time.h>
 FILE *fp;
main (argc, argv)
int argc;
char *argv[];
         start_str[25], stop_str[25];
 time_t st, et;
 struct tm start_t, end_t;
 double sec;
        get_line(), end_f;
 fp = fopen(argv[1], "r");
 while (1)
   end_f = get_line(start_str);
   if (end_f == EOF)
    break;
   end_f = get_line(stop_str);
   (void) strptime (start_str, "%a %h %e %T %Y", &start t);
   (void) strptime (stop_str, "%A %B %d %T %Y", &end_t);
  st = timelocal(&start t);
  et = timelocal(&end t);
  printf ("\t%4.2f\n", (et - st)/3600.0);
  /* (void) difftime(st, et); */
}
int get line (line)
 char
            *line;
 char ch;
 int i = 0;
 ch = fgetc (fp);
 while (!(ch == \n' || ch == EOF))
  i++;
  *line++ = ch;
  ch = fgetc (fp);
 *line = \0';
 if (ch == EOF)
  return EOF;
 else
  return i;
}
```

RSEED CPL

```
##This script extracts the the initial seed values at the time of
#convergence from FILE1 and places them in FILE2.
#

if ($3 != "") then
    awk -f ~rsimpson/script_dir/rand_awk Rmax=$1 $2 >tr_tmp_101
    tr -d ',' < tr_tmp_101 > $3
    rm tr_tmp_101

else if ($2 != "") then
    awk -f ~rsimpson/script_dir/rand_awk Rmax=$1 $2 >tr_tmp_101
    tr -d ',' < tr_tmp_101
    rm tr_tmp_101
    rm tr_tmp_101

else
    echo "usage: rnd 0-9 run_file [seed_file]"
endif
```

CALG_CPL

```
#This script extracts the performance statistics for individual
#algorithms from the FILE1 -containing alternate data for each.
#The results are placed in individual temporary files before
#being concatenated and stored in FILE2.
#All temporary files are then erased.
set fid=temp
if (($1 != "") && ($2 != "")) then
 set Cx=1
 if (\$3 != "") set fid = \$3
 while (Cx \le 1)
  awk -f ~rsimpson/script dir/alg awk idx=$Cx algs=$1 $2 >$fid{ alg}$Cx
  end
 if ($3 != "") then
  cat /dev/null >$3
  foreach file ($fid{_alg}*)
    cat $file >>$3
  end
 else
  foreach file ($fid{_alg}*)
   cat $file
  end
  rm $fid{ alg}*
 endif
else
   echo "usage: cg 0-9 stat_file [ssort_file]"
endif
```

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CONVERT_CPL

```
#!/bin/csh
#
# This script extracts the initial simulation trace file (FILE1)
# the results at the time of convergence from FILE1 and places them
# in FILE2.
#
if ($2 != "") then
    awk -f ~rsimpson/script_dir/stat_awk $1 >tr_tmp_101
    tr -s '' '\11' <tr_tmp_101 > $2
    rm tr_tmp_101
else if ($1 != "") then
    awk -f ~rsimpson/script_dir/stat_awk $1 >tr_tmp_101
    tr -s '' '\11' < tr_tmp_101
    rm tr_tmp_101
else
    echo "usage: cv run_file [stat_file]"
endif</pre>
```

CRSHEET_CPL

```
# This script creates the data to be used in the spreadsheet.
if ($2 != "") then
  cv $2 >tr tmpst 101
  gt $2 >tr_tmpti 101
  rnd $1 $2 >tr_tmpsd_101
  paste \ tr\_tmpst\_101 \ tr\_tmpti\_101 \ tr\_tmpsd\_101 > tr\_tmpRS\_101
  if ($3 != "") then
    tr -s * ' '
                      '<tr_tmpRS_101 >$3
  else
    tr -s '
                      '<tr_tmpRS_101
  endif
  rm tr_tmp*_101
 echo "usage: go 0-9 run_file [stat_file]"
endif
```

EXCEL CPL

```
#!/bin/csh

#

#

if (($1 != "") && ($2 != "")) then
    go $2 $1 > x_tmp_101
    cg $2 x_tmp_101
    rm x_tmp_101

else
    echo "usage: xcel run_file [No_Alg]"
endif
```

GET TIME.CPL

```
#This script will extract the start and stop times from
#each simulation run - stored in FILE1, and store them
#in FILE2.
set V = $2
if ($2 == "") then
 set V = temp_t123.dat
endif
if ($1 != "") then
 awk -f ~rsimpson/script_dir/rtime_awk $1 >$V
#The difference in time (hrs) of each run is computed
#and stored along with start and stop times in file3.
 if ($3 != "") then
    ~rsimpson/script_dir/calctime.exe $V >$3
    ~rsimpson/script dir/calctime.exe $V
    if ($V == "temp_t123.dat") then
      rm $V
    endif
 endif
else
 echo "usage: gt run_file [[tim_file] [hrs_file]]"
endif
```

C.3 ANSI C IMPLEMENTATION

This section presents the final and complete source code for the simulation model described in this study:

$Model_params.h$

	4
#define	SUCCESS 1
#define	FAILED 0
#define	ASCEND 1
#define	DESCEND 0
#define	EMPTY -1
#define	BOOLEAN int
#define	TRUE 1
#define	FALSE 0
#define	QUEUES_FLUSHED 1
#define	DATA_CONVERGED 2
#define	NQUANTA 10
#define	N_SYS_PROCS 1
#define	SAMPLE T 10
#define	SMPBLK_SZ (SAMPLE_T + 1)
#define	TABLE_SIZE 4
#define	CSET_SIZE 8
,, 4010	COD1_012D 0
#define	CONTEXT_SWITCH 200
#define	QUANTUM 20000
#define	ONE_SECOND 1000000.0
#define	ONE_MINUTE (60 * ONE_SECOND)
#define	ONE_HOUR (60 * ONE_MINUTE)
#define	BAKUP_INTERVAL (8 * ONE_HOUR)
#define	DUMP_INTERVAL (10 * ONE_SECOND)
#define	STABLE STATE 3600
#define	TRACE_INTERVAL (5 * ONE HOUR)
#define	AVE PROC GROUP 1.0
#define	AVE_EXEC_TIME 1.0
#define	AVE_INST 1.0
#define	KC CALL 0
#define	RCV_AVG_TOUT -11
#define	EXEC_CALL -10
#define	PRB_TOUT -9
#define	CANCEL TMER -8
#define	RESET_TMER -7
#define	SRQ TOUT -6
#define	WAITP_TOUT -5
#define	LOW TOUT -4
#define	HIGH TOUT -3
#define	CREATE CALL -2
#define	HALT_CALL -1
#define	KC EMSG 1
#define	LCAL_PRB_MSG 2

```
#define
                 PSTOP_MSG
                                  3
#define
                 PRB_MSG
                                  4
#define
                 PRB RPLY MSG
                                  5
#define
                 PRC ARI MSG
                                  6
#define
                 WK_RQST_MSG
                                  7
                 WK_RPLY_MSG
#define
                                  8
#define
                 SND_PRC_MSG
#define
                 HIGH LD MSG 10
#define
                 NEW_LD MSG
                                11
#define
                 LOAD STATE
                                 12
#define
                 LOW_LD_MSG 14
#define
                 PINFO_MSG
#define
                 TIME EXIT
                                   200
#define
                 RX BYTE TIME
                                   1.0
                 TX_BYTE_TIME
#define
                                   1.0
#define
                 PROTOCOL TIME 1000.0
#define
                 TIME_SLICE
                                   50000.0
#define
                 ELAPSED TIME
                                   10000.0
#define
                 SYNC_TIME
                                    50000.0
#define
                 EXECUTION_TIME (AVE_EXEC_TIME * ONE_SECOND)
#define
                 USERtime
                                    1
#define
                 OStime
#define
                 TIME_MIG_PROC (PROTOCOL TIME * AVE INST)
#define
                 TIME_RCV_PROC (PROTOCOL_TIME * AVE_INST)
#define
                 MESS_SIZE
                               10000.0
#define
                 START NODE
                                      15
#define
                 CONVERGE FACTOR 0.02
#define
                 PAGE SIZE
                                     500000
#define
                 BATCH RUN
                                      FALSE
/*
#define
                 DEBUG DUMP
                                      FALSE
#define
                 MIN_PID
                                      008
#define
                 MAX PID
                                      1200
#define
                 DEBUG MSG
                                      TRUE
#define
                 TIMEOUT
                                     (EXECUTION_TIME / 5.0)
#define
                                     (EXECUTION TIME / 4.0)
                 GBAL TOUT
                                    (2*TX_BYTE_TIME+(PROTOCOL_TIME*
#define
                 ARR TIME
AVE INST))
                               (ARR_TIME * msg->distance * 2 + TIMEOUT)
#define
                  Delay Time
#define
                  context_switch_time_update(node_P, (int) (CONTEXT_SWITCH * AVE_INST),
USERtime)
```

```
These type declaration differe from the original through the
  introduction of a buffer number and an "event filename" per cpu.
 typedef struct processor_node PROCESSOR;
typedef struct process_entry USER_PROCESS;
typedef struct comm_packet MSG FRAME;
typedef struct data_frame GLBAL_FRAME;
typedef struct
     int actual load;
     int virtual load;
     int OSportion;
     int used;
    QTUM ENTRY;
typedef struct
     short int
                  node id;
     PROCESSOR
                  *node_ptr;
    ROUTE_ENTRY;
typedef struct
     int
                 node id,
                 hops;
    LINK_ENTRY;
typedef struct
     short int
               id,
                ld:
   LOAD_ENTRY;
typedef struct
     BOOLEAN
                      await sync;
     PROCESSOR
                      *node_adr;
   NODE_ENTRY;
```

```
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```

```
struct process_entry
  int
                       upid;
  int
                       residency;
  int
                       n probes;
  double
                      exec_time;
  int
                      exec_here_time;
  int
                       residency_time;
  double
                      exist time;
  int
                      orig_mc,
                       preferred mc;
  BOOLEAN
                      schedulable,
                       blocked,
                       finished,
                       mcs_probed[6];
  double
                      timeout;
  int
                       itype;
  USER_PROCESS
                      *nxt_proc_ptr;
 };
 struct comm_packet
  double
                      time_stamp;
  int
                      sender_id,
                      dest_id,
                      distance,
                      total bytes,
                      type,
                      service_no;
   double
                      ldvalue;
  BOOLEAN
                      ack_reply;
   USER_PROCESS *mess_data;
  MSG FRAME
                     *next_pack_ptr;
};
struct data_frame
  int
                       size,
                        type,
                        sno;
  USER_PROCESS
                     *prc;
  double
};
typedef struct
      double
                        mintime;
                         itype;
      USER_PROCESS *minP;
    TMER_QUEUE;
typedef struct
      MSG_FRAME
                      *head,
                       *tail;
    MESS_QUEUE;
```

```
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```

```
struct processor_node
  QTUM_ENTRY
                       quanta[NQUANTA];
  int
                       quanta_idx,
                        OSoverhead;
  short int
                      node id;
  unsigned short
                    xsubi[3];
  LINK_ENTRY
                     *links;
 LOAD_ENTRY
                      cSet[CSET_SIZE];
 int
               load,
               threshold,
               n_tx, n_rx,
               n_migrates,
               n_immigs,
               n deaths,
               n_local_procs,
               n_active_local_procs,
               n_virtual_procs;
 float
              wt;
 BOOLEAN high_t_set,
              low_t_set,
             pwait_set,
              rplimit;
 double
              sys_real_time,
              next_elapsed_second,
              last_perf_dump_time;
 ROUTE ENTRY
                        route_table[TABLE_SIZE];
 double
                 nxt_arr_time,
                 cumul_exist_time;
 USER_PROCESS
                        *curr_process,
                        *firstjb ptr,
                        *lastjb_ptr;
 MESS_QUEUE
                       queue;
 TMER_QUEUE
                       timq;
 short int
                    buffer_no;
 char
                    event_fname[16];
FILE
                      *job_file,
                      *trace_file;
```

};

```
This version of initb_module on SparcStation One has
                                                                         */
     been modified inorder to enable meshes greater than
                                                                         */
     4x4 to be modelled. This was achieved by the removal
                                                                         */
     of a recursive descent search in building the 'shortest
     path route table.
                                                                        */
                                                                        */
     Date of Modification: 27th March 1991
                                                                        */
#include <stdio.h>
#include <stdlib.h>
#include <strings.h>
#include <math.h>
#include "model_params.h"
#include "model types.h"
#include "globvar.h"
extern FILE
                *trace,
               *seedfile,
                *monitor file,
                *configfile,
                *paramfile;
typedef struct stack_entry
 short int
                     node id,
                     distance;
 BOOLEAN
                    visited;
 ROUTE ENTRY
                     aroute;
 struct stack_entry
                   *next;
} STACK_ENTRY;
NODE ENTRY
                      *processor_table;
int
              *VTable;
               VTotal;
float
int
               mesh length;
extern int
               total_pe;
STACK_ENTRY
                       *main_stak,
                       *dum stak;
static int
               gbal stak[36],
                gbal_t;
BOOLEAN
               Key_In_List (int [],int, int);
               free_qspace(STACK_ENTRY *),
void
               push(STACK_ENTRY **, STACK_ENTRY *);
BOOLEAN
               Balanced Route (PROCESSOR *, int);
STACK_ENTRY *Create_ItemQ(int, BOOLEAN,int, PROCESSOR *),
                  *pop(STACK ENTRY **);
```

```
Ave_Call
                                                                           */
int Ave_Call (PROCESSOR *node_P)
  int
          idx,
          sum = 0,
          nbors = 0;
  for (idx = 0; idx < TABLE_SIZE; idx++)
    if (node_P->route_table[idx].node_id == EMPTY)
      continue;
    else
       sum += VTable[node_P->route_table[idx].node_id];
       nbors++;
    };
  return ((sum)/nbors);
     dump_stack
void dump_stack(STACK_ENTRY *top)
while (top != NULL)
 printf ("\nP\%d\tN\%d\tkm = \%d\t", top->node_id,
      top->aroute.node_id, top->distance);
 if (top->visited)
   printf ("TRUE");
   printf ("FALSE");
 top = top->next;
printf ("\n\n");
```

```
Reachable
                                                                                  */
BOOLEAN Reachable (int pid, int sn, int dn, int km)
   int
         tp = 0, stak[36];
   int
         i, nd, cpy;
   stak[tp++] = -(pid + 1);
   cpy = km;
   do{
     stak[tp++] = -(sn + 1);
     for (i = 0; i < TABLE\_SIZE; i++)
      nd = processor_table[sn].node_adr->route_table[i].node_id;
      if (nd == EMPTY)
        continue;
      if (Key_In_List(stak, tp, nd + 1));
        stak[tp++] = nd + 1;
      km --;
 /*
      printf ("\nReachability Stack Dump km =%d\n", km);
      for (i = 0; i < tp; i++)
        printf ("%d\n", stak[i]);
      if (Key_In_List(stak, tp, dn + 1) \&\& (km == 1))
       return TRUE;
      if (km \ll 1)
       km++;
       while (stak[--tp] >= 0);
     if (tp > 1)
       while ((tp != 1) \&\& (stak[--tp] < 0)) km++;
       sn = stak[tp] - 1;
   }while (tp != 1);
   return FALSE;
    Key_In_List
BOOLEAN Key_In_List (int stk[], int t, int key)
 int i;
 for (i = 0; i < t; i++)
 if (key == (int) fabs((double)stk[i]))
    return TRUE;
 return (FALSE);
```

```
Update_Link
                                                                       */
BOOLEAN Update_Link (PROCESSOR *src, int nbor,STACK_ENTRY *top)
  int
          dn;
  dn = top->aroute.node id;
  if (dn != src->node id)
    if (Reachable (src->node_id, nbor, dn, top->distance))
     if (src->links[dn].node\ id < 0)
      VTable[nbor]+= 1;
      src->links[dn].node_id = nbor;
      src->links[dn].hops = top->distance;
     free_qspace(top);
     return (TRUE);
    };
  push (&main_stak, top);
  return (FALSE);
   write_intro_script
                                                                      */
                                                                      */
void write_intro_script (void)
  char
          net_script[512];
  int
         idx;
  for (idx = 0; idx < 512; idx++)
  if (feof(configfile))
    break;
  if ((net_script[idx] = getc (configfile)) == '!')
    break;
 net_script[idx] = '\0';
 idx = 0;
 while (net\_script[idx] != '\0')
  putc (net_script[idx], monitor_file);
  idx++;
 fscanf (configfile, "%d", &total_pe);
```

```
Create_proc_node
                                                                         */
                                                                         */
int Create_proc_node (int node_id,NODE_ENTRY processor_table[])
{ /* implementation */
 PROCESSOR *new_node;
 new_node = (PROCESSOR *) malloc (sizeof(PROCESSOR));
 if (new_node == NULL)
   return FAILED;
 else
  { int
         idx;
   /* Construct data node */
   processor_table[node_id].node_adr = new_node;
   new_node->node_id
                          = node_id;
   new_node->links = (LINK_ENTRY *)
             calloc (total_pe, sizeof (LINK_ENTRY));
   for (idx = 0; idx < total_pe; idx++)
     new_node->links[idx].node_id
                                       = EMPTY;
     new_node->links[idx].hops
                                      = total_pe;
   for (idx = 0; idx < 4; idx++)
     new_node->route_table[idx].node_ptr = NULL;
    new_node->route_table[idx].node_id = EMPTY;
   return SUCCESS;
```

```
auto_link_creation
                                                                               */
void auto_link_creation (FILE *configfile, int mesh_len)
 int node_no,
    row_idx, col_idx;
 node no = 0;
 for (row_idx = 0; row_idx < mesh_len; row_idx++)
   for (col_idx = 0; col_idx < mesh_len; col_idx++)
    fprintf (configfile, "%d\n", node_no);
    if (row_idx < (mesh_len -1))
      fprintf (configfile, "%d %d\n",
           node_no, node_no + mesh len);
     if (row idx > 0)
       fprintf (configfile, "%d %d\n",
             node_no, node_no - mesh len);
    }
    else
      fprintf (configfile, "%d %d\n",
           node_no, node_no - mesh_len);
    if (col_idx < (mesh_len -1))
     fprintf (configfile, "%d %d\n",
           node_no, node_no + 1);
     if (col_idx > 0)
       fprintf (configfile, "%d %d\n",
             node_no, node_no - 1);
    }
    else
     fprintf (configfile, "%d %d\n",
          node_no, node_no - 1);
    node_no++;
 fprintf (configfile, "%d\n", node_no);
```

```
manual_link_create
                                                                                   */
void manual_link_create (FILE *configfile, int maxmcs)
 int link no,
     node_id;
 for (node_id = 0; node_id < maxmcs;)</pre>
  fprintf (configfile, "%d\n", node id);
  printf ("\nNode%d is:", node_id);
  for (;;)
    printf ("connected to ->");
    scanf ("%d", &link_no);
    if ((link\_no < 0) \parallel (link\_no == node\_id)
       || (link_no > maxmcs))
      break;
    fprintf (configfile, "%d %d\n", node_id, link_no);
  node id++;
fprintf (configfile, "%d\n", node_id);
     square mesh
                                                                                  */
    square_mesh (int total_pe)
double value;
 value = sqrt((double) total_pe);
 return (int) (floor (value));
```

```
create_config file
                                                                                */
void create_config_file (BOOLEAN batch run)
 char response[3],
       net_description[512];
       pe_max, prev_total = 0;
 int
 if (!batch run)
  printf ("\nDo you wish to create a new configuration file?");
  scanf ("%s", response);
  if (strcmp(response, "no\0"))
     configfile = fopen ("config.dat", "w");
     for (;;)
      printf ("\n\nEnter the total number of PE's ");
      scanf ("%d", &pe max);
      if (pe max \leq 1)
       break;
      total pe = pe max;
      printf ("\nPROVIDE A BRIEF DESCRIPTION OF THE NETWORK TOPOLOGY AND
THE(n''):
      printf ("\nPROCESSOR ARRANGEMENT YOU INTEND TO SIMULATE...\n");
      printf ("\nTHE TEXT MUST BE TERMINATED BY AN EXCLAMATION MARK\n\n");
      {int idx;
       for ( idx = 0; idx < 512; idx++)
        if ((net_description[idx] = getc(stdin)) == '!')
          break:
       net description[++idx] = \0';
      fprintf (configfile, "%s\n", net_description);
      fprintf (configfile, "%d\n", total_pe);
     mesh_length = square mesh (total pe);
     if (total pe == prev total)
       printf ("\nDo you wish to use the preceding config. job stream?");
       scanf ("%s", response);
       fprintf (configfile, "%d\n",
            (int) ((strcmp(response, "no\0")));
     }else prev_total = total_pe;
     fprintf (configfile, "%d\n", mesh length);
     if (mesh length * mesh length == total pe)
       printf ("\nAutomatic point-to-point connection? ");
       scanf ("%s", response);
       if (strcmp(response, "no\0"))
         auto link creation (configfile, mesh_length);
       else manual_link_create (configfile, total_pe);
     }else manual link create (configfile, total_pe);
    fclose(configfile);
```

```
create_params_file
                                                                                 */
void create_params_file (BOOLEAN batch_run)
 void put_lba_parameters(int);
 char response[3];
 float load value;
int total_jobs,
      lba_index;
 if (!batch run)
  printf ("\nDo you wish to create a new PARAMETER file?");
  scanf ("%s", response);
  if (strcmp(response, "no\0"))
    paramfile = fopen ("pramfile.dat", "w");
    for (;;)
     printf ("\nEnter the LOAD value ");
     scanf ("%f", &load_value);
     if (load_value <= 0.0)
       break;
     fprintf (paramfile, "%f\n", load_value);
     printf ("\nEnter the total number of jobs: ");
     scanf ("%d", &total_jobs);
     if (total jobs < 1)
         break;
     fprintf (paramfile, "%d\n", total_jobs);
     printf ("\nWill this be a new job_stream? ");
     scanf ("%s", response);
     fprintf (paramfile, "%d\n",
          (int) ((strcmp(response, "no\0"))):
     printf ("\nSelect the required LBA: ");
     scanf ("%d", &lba_index);
     fprintf (paramfile, "%d\n", lba_index);
     put_lba_parameters(lba_index);
   fclose(paramfile);
```

```
put lba parameters
/*
                                                                                 */
void put_lba_parameters (int lba idx)
{ float threshold;
 int
        probe limit;
 probe_limit = 2 * (mesh_length - 1);
 switch (lba_idx)
  case 0: /* no load balancing */
       threshold = (float) 1000;
       fprintf (paramfile, "%f\n", threshold);
       break;
  case 1: /* Random
  case 3: /* Reverse
  case 4: /* Threshold Broadcast */
  case 9: /* Global Average - Sender */
  case 10:/* Global Average - Receiver */
       printf ("\nEnter the THRESHOLD level: ");
       scanf ("%f", &threshold);
       fprintf (paramfile, "%f\n", threshold);
       break;
  case 5: /* Threshold Neighbour */
  case 6: /* Adaptive Threshold */
  case 7: /* Adaptive Reverse */
  case 8: /* Global Average Nbor */
       printf ("\nEnter the THRESHOLD level: ");
       scanf ("%f", &threshold);
       printf ("\nEnter the Maximum Distance: ");
       scanf ("%d", &probe limit);
       fprintf (paramfile, "%f\n", threshold);
       break;
 case 2: /* Threshold Policy */
       printf ("\nEnter the THRESHOLD level: ");
       scanf ("%f", &threshold);
       fprintf (paramfile, "%f\n", threshold);
       printf ("\nEnter the PROBE LIMIT: ");
      scanf ("%d", &probe limit);
 default:break;
fprintf (paramfile, "%d\n", probe_limit);
```

```
Build_Network
                                                                                */
void Build_Network(void)
 int
            Create_proc_node(int, NODE_ENTRY[]);
 void
            Build_route_table(NODE_ENTRY []);
 int
            i,idx, lidx,
           node_id, link_id, prev_id;
 PROCESSOR
                  *node ptr;
 for (idx = 0; idx < total_pe; idx++)
   if (!Create_proc_node(idx, processor_table))
    break;
 fscanf (configfile, "%d", &node_id);
 idx = 0:
 while ((idx < total_pe) && (!feof(configfile)))
    lidx = 0:
    node_ptr = processor_table[idx].node_adr;
    prev id = node id;
    fscanf (configfile, "%d", &node id);
    while (node id == prev id)
     fscanf (configfile, "%d", &link id);
     node_ptr->route_table[lidx].node_id = link_id;
     node_ptr->route_table[lidx].node_ptr
       = processor_table[link_id].node_adr;
     node_ptr->links[link_id].node_id = link_id;
     node_ptr->links[link_id].hops = 1;
    lidx++;
     fscanf (configfile, "%d", &node_id);
   idx++;
 VTable = (int *) calloc (total_pe, sizeof (int));
 Build_route_table(processor_table);
 for (i = 0; i < total pe; i++)
   VTotal += VTable[i];
#ifdef DEBUG MSG
  dump_tables (processor_table);
#endif
 for (i = 0; i < total_pe; i++)
  processor_table[i].node_adr->wt = VTable[i]/ VTotal;
  VTable[i] = 1;
 VTotal = total pe;
```

```
Build_route_table
                                                                           */
void Build_route_table(NODE_ENTRY processor_table[])
 int idx;
 void traverse_net(PROCESSOR *);
    for (idx = 0; idx < total_pe; idx++)
       VTable[idx] = 1;
    idx = START_NODE;
    do
    {
     main_stak = NULL;
     traverse_net(processor_table[idx].node_adr);
     free_qspace(main stak);
     idx = (idx + 1) \% total_pe;
    } while (idx != START NODE);
    VTotal = total pe;
}
     stack links
                                                                          */
BOOLEAN stack_links (PROCESSOR *p, int km)
short int
             node_id, i;
STACK_ENTRY *stack_item;
BOOLEAN
                 stacked = FALSE,
          already_stacked (int);
stack_item = Create_ItemQ (km, TRUE, EMPTY, p);
push (&main_stak, stack_item);
for (i = 0; i < TABLE\_SIZE; i++)
 node_id = p->route_table[i].node_id;
 if ((node id != EMPTY) && (!already stacked(node id)))
  push (&main_stak, Create_ItemQ (km +1, FALSE, i, p));
    stacked = TRUE;
 return stacked;
```

```
Incomplete_Links
                                                           */
BOOLEAN Incomplete_Link (LINK_ENTRY ltab[])
  int
       count = 0;
  for (i = 0; i < total_pe; i++)
    if (ltab[i].node id < 0)
     count++;
  if (count == 1)
   return FALSE:
  else
   return TRUE;
Find_Links
BOOLEAN Find_Links (LINK_ENTRY ltab[],int nbor,int km)
 int
       i;
 BOOLEAN found = FALSE, stacked;
 for (i = 0; i < total_pe; i++)
   if ((ltab[i].node_id == nbor)
     && (ltab[i].hops == km))
     stacked = stack_links(processor_table[i].node_adr, km);
     found = stacked || found;
 return stacked;
          */
    Push Ltable
BOOLEAN Push_Ltable (PROCESSOR *src)
{int
         success = FALSE,
         nbour id,
         stak[4];
short int
         i, t = 0;
STACK_ENTRY *top;
top = pop(&main stak);
while (top ->distance != 0)
 push (&main_stak, top);
 for (i = 0; i < TABLE\_SIZE; i++)
  nbour_id = src->route_table[i].node_id;
  if (nbour id == EMPTY)
   continue;
  else
   if (Balanced Route (src, nbour_id))
```

```
success = TRUE:
       break;
     }
     else
      if (!Key_In_List(stak, t, nbour_id))
        stak[t++] = nbour_id;
    while (t > 0)
      nbour_id = stak[--t];
      top = pop(&main_stak);
      if (Update_Link(src, nbour_id, top))
        success = TRUE;
        break;
    top = pop(&main_stak);
 push (&main_stak, top);
 return success;
     traverse net
                                                                              */
                                                                              */
void traverse_net (PROCESSOR *src)
{
int
                 nbour_id, i,
                 distance;
 BOOLEAN
                 found, stacked = FALSE,
                 Incomplete_Link (LINK_ENTRY[]);
distance = 1;
 push (&main_stak, Create_ItemQ(0, TRUE, EMPTY, src));
while (Incomplete_Link(src->links))
  for (i = 0; i < TABLE\_SIZE; i++)
   nbour_id = src->route_table[i].node_id;
   if (nbour id == EMPTY)
    continue;
   else
    Find_Links (src->links, nbour_id, distance);
    Balanced Route (src, nbour id);
  if (Push_Ltable(src))
     distance++;
  else distance--;
free_qspace(pop(&main_stak));
```

```
Balanced_Route
                                                                           */
BOOLEAN Balanced_Route (PROCESSOR *src, int nbor)
                succeed = TRUE;
 BOOLEAN
 int
          dn;
 STACK_ENTRY *top;
 top = pop(&main_stak);
 dn = top->aroute.node id;
 if (dn == src->node_id)
  push(&main_stak, top);
 else
  if (src->links[dn].node_id < 0)</pre>
    if (VTable[nbor] < Ave_Call(src))</pre>
      succeed = Update_Link(src, nbor, top);
       push (&main_stak, top);
       succeed = FALSE;
   }
   else
    free_qspace(top);
    return (Balanced_Route (src, nbor));
return succeed;
     already_stacked
                                                                           */
BOOLEAN already_stacked (int id)
  struct stack entry *dumq;
  dumq = main stak;
  while (dumq != NULL)
    if (dumq->aroute.node_id == id)
     return (TRUE);
     dumq = dumq->next;
  return (FALSE);
```

```
free_qspace
                                                                */
void free_qspace(struct stack_entry *q_entry)
 struct stack_entry *sublist ptr;
 while (q_entry != NULL)
   sublist_ptr = q_entry->next;
   free ((char *) q_entry);
   q_entry = sublist_ptr;
    push
                                                               */
void push(STACK_ENTRY **stak, STACK_ENTRY *entry)
 if (entry != NULL)
   if (stak == NULL)
   *stak = entry;
    entry->next = *stak;
    *stak
            = entry;
                                                               */
    pop
                                                               */
STACK_ENTRY *pop(STACK_ENTRY **stak)
 STACK_ENTRY *top;
 if (*stak == NULL)
   return (NULL);
 else
  {
           = *stak;
   top
   *stak
          = top->next;
   top->next = NULL;
   return (top);
```

```
Create_ItemQ
                                                                         */
struct stack_entry *Create_ItemQ(int km,
                    BOOLEAN visited,
                    int route_idx,
                    PROCESSOR *p)
  struct stack_entry *q_entry;
  q_entry = (struct stack_entry *) malloc (sizeof(struct stack_entry));
  if (q_entry == NULL)
    return (NULL);
  else
    q_entry->node_id
                         = p->node id;
    q entry->distance
                           = km;
    q entry->visited = visited;
    if (route idx != EMPTY)
     q_entry->aroute = p->route_table[route_idx];
      q_entry->aroute.node_id = p->node id;
      q_entry->aroute.node_ptr = p;
    q_entry->next = NULL;
    return (q_entry);
     Dump MyLink
                                                                         */
/<del>*</del> _____ */
    Dump_MyLink(LINK_ENTRY Itab[], int id)
 int
           j;
  fprintf (monitor_file, "\n\nProcessor%d PRIMARY route\n", id);
  for (j=0; j < total pe; j++)
   fprintf (monitor_file,
         "\ndest.:%d via:node%d distance:%d nodes",
         j, ltab[j].node id,
         ltab[j].hops);
  */
  fprintf (monitor_file,
       "\n\n***%8d ROUTES via this node ***\n",
       VTable[id]);
  return (0);
```

```
dump tables
/*
                                                                              */
/* ==
     dump_tables(NODE_ENTRY processor_table[])
int
 {int
 PROCESSOR
                  *node_ptr;
 for (i = 0; i < total pe; i++)
  node_ptr = processor_table[i].node_adr;
  Dump_MyLink (node_ptr->links, i);
   printf ("\nProcessor%d SECONDARY route\n", node_ptr->node_id);
  for (j=0; j < total_pe; j++)
   fprintf (monitor_file, "\ndest.:%d via:node%d distance:%d nodes\n",
        j, node_ptr->links[j][2].node_id,
        node_ptr->links[j][2].hops);
 return (0);
/**** only used when debugging the routing strategy ***
main (void)
 int msh_size;
 create config file (0);
 create_params_file (0);
 configfile = fopen ("config.dat", "r");
 write_intro_script();
 processor_table = (NODE_ENTRY *)
            calloc (total_pe, sizeof(NODE_ENTRY));
 fscanf (configfile, "%d", &msh_size);
 Build_Network();
 dump_tables (processor_table);
 fclose (configfile);
         THEEND
```

```
This version of kernal_module on SparcStation One has
     been modified so that trace dumps can be performed over
/*
/*
     larger and longer time intervals.
                                                                  */
                                                                  */
     Date of Modification: 8th April 1991
                                                                 */
#include <stdio.h>
#include <stdlib.h>
#include <strings.h>
#include "model_params.h"
#include "model types.h"
#include "globvar.h"
/***********************
     CONSTANT, GLOBAL, & LITERAL
       DEFINITIONS.
/**********************
/* ** ANSI restriction **
 jmp_buf
                 ErrLabel;
int
             Sec = 0, Del = 0,
             FinL, MaxL:
int
             process_count;
double
            global_clock,
            global_dump time,
            stop time;
unsigned short
                 n_reached;
char
             **io buffer blk;
int
            Virtual_Sender (PROCESSOR *),
           Virtual_Receiver(PROCESSOR *);
BOOLEAN
                 SkipRun,
                 finish:
extern NODE ENTRY
                    *processor table;
extern float
              threshold;
extern FILE
               *monitor file;
extern float
              Mprofile;
extern int
              probe limit,
              last_t,
              total pe;
extern GLBAL FRAME
                      *dblock:
extern int
                       Iba index:
```

```
A CALLEGE WHEN THE
```

extern void time_update(PROCESSOR *, int, int);

send_probe(PROCESSOR *, USER_PROCESS *),

new_load_msg (PROCESSOR *, MSG_FRAME *),

probe_reply_msg (PROCESSOR *, MSG_FRAME *),

work_reply_msg (PROCESSOR *, MSG_FRAME *),

lcal_reply_msg (PROCESSOR *, MSG_FRAME *),

Gbal_reply_msg (PROCESSOR *, MSG_FRAME *),

transfer_msg (PROCESSOR *, struct comm_packet *),

process_msg (PROCESSOR *, MSG_FRAME *),

exit_msg (PROCESSOR *, MSG_FRAME *),

exit_msg (PROCESSOR *, MSG_FRAME *),

exit_proc(PROCESSOR *, MSG_FRAME *);

extern BOOLEAN

probe_msg (PROCESSOR *, MSG_FRAME *),
work_request_msg (PROCESSOR *, MSG_FRAME *);

extern GLBAL_FRAME *pack_data(USER_PROCESS *, int, int);

extern void

close_eventfiles(void),

ave_compute_or_display(BOOLEAN),

final_compute(void), psw_backup();

extern void

Select_next_timeout(PROCESSOR *);

extern BOOLEAN simulation_complete(void),

Bcast_Cset(PROCESSOR *, int, int, int),

int_divisible(double, float);

extern char

*Get_Time(void);

extern int

init_eventfile(PROCESSOR *, short int);

extern MSG_FRAME

*receive(PROCESSOR *);

extern MSG_FRAME

*msg_packet (PROCESSOR *, int);

extern BOOLEAN

page_faulting(PROCESSOR *);
migrate(PROCESSOR *, USER_PROCESS *, int);

extern void extern void

broadcast(PROCESSOR *);

extern int

bcast_nbor (PROCESSOR *, int, int);

extern void

restart_clock(PROCESSOR *, int, double, USER_PROCESS *);

extern void

Update_Cset (PROCESSOR *, MSG FRAME *, int);

extern int

Cset_Load_Avg (PROCESSOR *);

```
SIMULATE_NETWORK
                                                              */
void
     simulate_network(void)
{ PROCESSOR
                *node P;
BOOLEAN
                 run_kernel(PROCESSOR []);
void
             initialise_nodes(void),
             reset_reached(void);
int
             exit_code;
initialise_nodes ();
reset_reached();
"TIME", "VAR.", "LOAD", "DIFF.", "RTime", "%Conv", "MIG.", "Tx", "%JOB");
"---", "---", "----", "----", "----", "----", "----"):
SkipRun = FALSE;
finish = simulation_complete();
while (!finish)
 { int idx;
 for (idx = 0; idx < total_pe; idx++)
   if (finish = !run_kernel (processor_table[idx].node adr))
     break;
 if (global_clock < stop_time)
  global_clock = global_clock + ELAPSED TIME;
 if ((n reached >= total pe)
   && (global_clock >= stop_time))
   reset reached();
  if (global_clock >= global_dump_time + DUMP_INTERVAL)
   if (int_divisible(global_clock, TRACE_INTERVAL) == TRUE)
     ave compute or display(TRUE);
     fflush (monitor_file);
   global_dump_time = global_clock;
 exit code = simulation complete();
 finish = finish || exit_code || SkipRun;
```

```
/* ** ANSI restriction ** *
  if (setjmp(ErrLabel) > 0)
      fprintf (monitor_file,
      "\n\n::M A L L O C ****** FAILED");
     fprintf (monitor_file, "in msg_packet()\n");
     fprintf (monitor_file, "Process No.%d\n", process_count);
     finish = TRUE;
** ** */
 } /* END SIMULATION */
 final_compute();
 Sec = (int) (global_clock / ONE_SECOND);
 fprintf (monitor_file, "\n *** TERMINATING CONDITION :");
 if (exit_code == DATA_CONVERGED)
   fprintf (monitor_file,
        " CONVERGENCE @%2.2f%%***",
           100 * CONVERGE_FACTOR);
 else
   fprintf (monitor_file, " JOB STREAM EXHAUSTED ***");
 fprintf (monitor file,
        " CR/Sec: %2.2f, DEL/Sec: %2.2f\n",
         (float) (process_count/Sec),
         (float) (Del/Sec));
 fprintf (monitor file,
        "\nFinal Th: %d, Max Th: %d\n", FinL, MaxL);
 fprintf (monitor_file, "\n\nSIMULATION COMPLETED ON %s\n", Get_Time());
/* Dump_RTime(); */
close eventfiles();
} /* END simulate network() */
```

```
/*********************
     INITIALISE_NODES
                                                  */
/***************
     initialise nodes (void)
 extern float Mean_load,
             sigma_load,
             sigma_var,
             sigma rtime,
             sigma_migrate,
             sigma_Tx;
 extern int
                sigma_diff;
 short int
                idx;
 char
                trace fname[16];
 PROCESSOR *node_P;
 global_clock = global_dump_time = 0.0;
 stop time = 0.0;
 process_count = 0; Del = 0;
 Mean_load = 0.0;
 sigma_load = 0.0, sigma_var
 sigma_rtime = 0.0, sigma_Tx
                               = 0.0,
 sigma migrate = 0.0;
 sigma_diff = 0;
 last_t = -1;
 Mprofile
            = 0.0; FinL = MaxL = (int) threshold;
 io_buffer_blk = (char **) calloc (total pe, sizeof(char *));
 for (idx = 0; idx < total pe; idx++)
  node_P = processor_table[idx].node_adr;
  time_update (node_P, (int) CONTEXT_SWITCH, OStime);
  *(io buffer blk + idx) = calloc (BUFSIZ, sizeof(char));
  sprintf (node_P->event_fname, "jobs%d_pg%d",
       idx, node_P->buffer_no);
  (void) init_eventfile (node_P, idx);
  node_P->wt = GBAL_TOUT;
#ifdef DEBUG DUMP
   sprintf (trace_fname,"trace%d_run%d",
        node_P->node_id, run_idx);
   node_P->trace_file = fopen (trace_fname, "w");
#endif
}
```

```
reset_reached
                                                        */
void reset_reached(void)
{int i;
 n_reached = 0;
for (i = 0; i < total_pe; i++)
  processor_table[i].await_sync = FALSE;
stop_time = stop_time + SYNC_TIME;
/*
    kernal_msg_handler
                                                        */
                                                        */
void kernal_rnsg_handler (PROCESSOR *node_P,MSG_FRAME *msg)
          kernal_call(PROCESSOR *, MSG_FRAME *),
void
          external_message(PROCESSOR *, MSG_FRAME *);
if (msg != NULL)
  if (msg->type == KC_CALL)
    kernal_call(node_P, msg);
  else
   external_message (node_P, msg);
  msg->mess_data = NULL;
  free ((char *) msg);
```

```
SERVICE_ROUTINE
                                                                 */
                                                                 */
void Service_routine (PROCESSOR *node_P)
{
           prb_timout_msg(PROCESSOR *, USER_PROCESS *),
void
          srq_timout_msg(PROCESSOR *, USER_PROCESS *),
           low_timeout_msg(PROCESSOR *);
USER_PROCESS
                  *proc_P,
                  *prev_P;
proc_P = node_P->firstjb_ptr;
prev_P = proc_P;
do
 if (proc_P == NULL)
   proc_P = node_P->firstjb_ptr;
 else
   if ((proc P->timeout)
     && (proc_P->timeout <= node_P->sys_real_time))
    if (proc_P->itype == SRQ_TOUT)
     srq_timout_msg (node_P, proc_P);
     if (proc_P->itype == PRB_TOUT)
       prb_timout_msg(node_P, proc_P);
   proc_P = proc_P->nxt_proc_ptr;
  if (proc_P == prev_P)
    break;
 } while (proc_P != prev_P);
```

```
/*
    TIMEOUT_HANDLER
BOOLEAN TimeOut_Handler(PROCESSOR *nodeP)
{
 void
        wrq_timout_msg(PROCESSOR *),
     hi_timeout_msg(PROCESSOR *),
       Avgld_tout_msg(PROCESSOR *),
     low_timeout_msg(PROCESSOR *);
if (nodeP->timq.itype)
  if (nodeP->timq.mintime <= nodeP->sys_real_time)
   switch (nodeP->timq.itype)
    case WAITP_TOUT: wrq_timout_msg(nodeP);
             nodeP->rplimit = 0;
             nodeP->pwait_set = 0;
             nodeP->n_virtual_procs = 0;
             break;
    case HIGH_TOUT : hi_timeout_msg(nodeP);
             break;
    case LOW_TOUT : low_timeout_msg(nodeP);
             break;
    case RCV_AVG_TOUT:
             Avgld_tout_msg(nodeP);
             break;
        default: Service_routine(nodeP);
    Select_next_timeout(nodeP);
   return TRUE;
   else
   if (nodeP->timq.minP!= NULL)
     Service_routine(nodeP);
return FALSE;
```

```
run_kernel
                                                                          */
/*****************
BOOLEAN run_kernel(PROCESSOR *node_P)
 USER PROCESS
                      *proc_P;
 void
             kernal_msg_handler(PROCESSOR *,MSG_FRAME *),
             schedule_process(PROCESSOR *);
 int
             nxt job,
             sender_initiated(PROCESSOR *, USER_PROCESS *);
 BOOLEAN
             suspend kernal,
             retry;
 MSG_FRAME
                     *msg_entry;
 if (node_P->sys_real_time < stop_time)
  msg entry
               = receive (node P);
  kernal_msg_handler(node_P, msg_entry);
  TimeOut Handler(node P);
  if (msg_entry == NULL)
    break;
 } while (1);
 suspend_kernal = (node_P->sys_real_time >= stop_time);
 if (!suspend_kernal)
  if (feof(node_P->job_file))
  if (!(page_faulting (node_P)))
      return FAILED;
  if (!feof(node P->job file))
    if (node P->sys real time >= node P->nxt arr time)
    { /* context switch; */
     proc_P = Add_Process(process_count, NULL, node_P);
     sender initiated (node P, proc P);
     fread ((char *)&nxt_job, sizeof(nxt_job), 1, node_P->job_file);
     fread ((char *)& (node_P->nxt_arr_time), sizeof(node_P->nxt_arr_time),
         node_P->job_file);
     process_count++;
    }; /* end outermost if: next event */
```

```
if \ (node\_P\text{-}>sys\_real\_time > node\_P\text{-}>wt) \\
     node_P->wt = node_P->sys_real_time + GBAL_TOUT;
      switch (lba index)
      {int
            xload;
        case 2:
        case 3:
        case 4:
       case 5: dblock = pack_data (NULL, LOAD_STATE, 4);
            xload = node_P->n_local_procs + node_P->n_virtual_procs;
            if (Bcast_Cset(node_P, xload, threshold, 0));
              Bcast_Cset(node_P, xload, threshold, 1);
            free ((char *) dblock);
            break;
       case 6:
       case 9: Virtual_Sender (node P);
           break;
       case 10:Virtual_Receiver (node_P);
           break;
       default:;
       };
     }
     schedule_process (node P);
#ifdef DEBUG_DUMP
   if (node_P->sys_real_time
      >= (node_P->last_perf_dump_time + DUMP_INTERVAL))
    performance_dump(node P);
    node_P->last_perf_dump_time = node_P->sys_real_time;
#endif
 }
 else
    if (!processor_table[node_P->node_id].await_sync)
     n_reached++;
     processor_table[node_P->node_id].await_sync = TRUE;
  };
  return SUCCESS;
```

```
/*
     SENDER_INITIATED
/*
int sender_initiated (PROCESSOR *node_P, USER_PROCESS *proc_P)
{ int
        xload,
        thresh_Bcast (PROCESSOR *, USER_PROCESS *,int),
        random_policy(PROCESSOR *, USER_PROCESS *),
        threshold_policy(PROCESSOR *, USER_PROCESS *);
 switch (lba_index)
   case 8: /* Global Symmetric -Creation/Deletion */
       return Virtual_Sender (node_P);
   case 11: /* Global Symmetric -Creation/Deletion */
       return Virtual_Receiver (node P);
   default:
    xload = node_P->n_local_procs;
    if (xload > threshold + N_SYS_PROCS)
    switch (lba_index)
   case 1: /* Random lb strategy */
       return (random_policy(node_P, proc_P));
   case 2: /* Threshold lb strategy */
      return (threshold_policy(node_P, proc_P));
   case 5: /* Threshold Neighbour*/
   case 4: /* Threshold Broadcast */
       return (thresh_Bcast(node_P, proc_P, xload));
   default:;
   };
 return (FAILED);
```

```
/*
     VIRTUAL_SENDER
/*******************
int Virtual_Sender (PROCESSOR *nodeP)
{
   int xload,
       global_avg_snd(PROCESSOR *,int),
       global_avg_rcv(PROCESSOR *),
       Rev_avg_snd(PROCESSOR *, int);
   xload = nodeP->n_local_procs + nodeP->n_virtual_procs;
   if (xload > nodeP->threshold + N_SYS_PROCS)
   switch (lba index)
   case 6: /* Adaptive Threshold */
        return Rev_avg_snd (nodeP, xload);
   case 8: /* Global Symmetric -Creation/Deleteion */
   case 9: /* Global Time Interval */
        return global_avg_snd (nodeP, xload);
   default:;
   }
  else
    if (xload <= nodeP->threshold)
    switch (lba_index)
        break;
    case 8: /* Global Symmetric */
       return global_avg_rcv (nodeP);
    default:;
    };
  return (FAILED);
}
```

```
VIRTUAL_RECEIVER
            ***************
int Virtual_Receiver (PROCESSOR *nodeP)
{ int xload,
       Rev_avg_snd(PROCESSOR *, int),
       Rev_avg_rcv (PROCESSOR *, int);
  xload = nodeP->n_local_procs + nodeP->n_virtual_procs;
  if (xload <= nodeP->threshold)
  switch (lba index)
   case 10: /* Global Receiver */
       return Rev_avg_rcv (nodeP, xload);
   default:;
  }
  else
   if (xload > nodeP->threshold + N_SYS_PROCS)
   switch (lba_index)
   case 10: /* Global Receiver */
       Rev_avg_snd (nodeP, xload);
   default:;
   };
  return (FAILED);
}
RECEIVER INITIATED
/*********************
int receiver_initiated (PROCESSOR *node P)
        xload,
{ int
         reverse_policy (PROCESSOR *, int);
 switch (lba_index)
  case 8: /* Global Symmetric - Creation/Deleteion */
       return Virtual_Sender (node_P);
  case 11: /* Global Symmetric - Creation/Deleteion */
       return Virtual_Receiver (node_P);
    xload = node P->n_local_procs + node_P->n_virtual_procs;
    if (xload <= threshold)
    switch (lba_index)
     case 3: /* Reverse lb strategy */
         return (reverse_policy(node_P, xload));
     default: ; /* Lb 0 - no load balancing
         Lb 3 - Reverse - receiver-initiated */
return (FAILED);
```

```
RANDOM POLICY
                                                       */
int random_policy (PROCESSOR *node_P, USER_PROCESS *proc_P)
{ int
          dest_id;
 dest_id = randmc_id(node P);
 proc_P = Remove_Process(node_P, proc_P->upid);
 if (proc P!= NULL)
   migrate (node_P,
      proc P,
      dest_id);
   node_P->n_local_procs--;
   node_P->n_active_local_procs--;
  return SUCCESS;
}
         THRESHOLD POLICY
                                                       */
int threshold_policy ( PROCESSOR *node_P,USER_PROCESS *proc_P)
 proc_P->schedulable = FALSE;
 node P->n active local procs--;
 send_probe(node_P, proc_P);
 return SUCCESS;
/***********************
   THRESH BCAST
                                                       */
                                                       */
int thresh_Bcast (PROCESSOR *nodeP, USER_PROCESS *procP,int xload)
  procP->schedulable = FALSE;
  nodeP->n_active_local_procs--;
  dblock = pack_data (procP, PRB_MSG, 4);
  nodeP->rplimit++;
  if (Bcast Cset(nodeP, xload, threshold, 0));
    bcast_nbor(nodeP,
        probe limit,
        xload);
  free ((char *)dblock);
  restart clock(nodeP, SRQ_TOUT,
        2 * probe_limit * ARR_TIME + TIMEOUT,
        procP);
  return SUCCESS;
}
```

```
REVERSE POLICY
/***************************
int reverse_policy (PROCESSOR *nodeP, int xload)
 if (!nodeP->pwait set)
    /* broadcast to all nodes, desire to receive work */
    dblock = pack_data (NULL, WK_RQST_MSG, 4);
     nodeP->rplimit = Bcast_Cset (nodeP, xload, threshold, 1);
     if (!nodeP->rplimit)
       nodeP->rplimit = bcast_nbor(nodeP,
                  probe limit.
                  xload):
    free ((char *)dblock);
    restart_clock(nodeP, WAITP TOUT,
           2 * probe_limit * ARR_TIME + TIMEOUT,
    nodeP->pwait_set = nodeP->rplimit;
    nodeP->low t set
 return SUCCESS;
    ADAPTIVE_REVERSE
/**********************
int Rev_avg_rcv (PROCESSOR *nodeP, int xload)
  if (nodeP->low\ t\ set == 0)
   dblock = pack data (NULL, LOW LD MSG, 4);
   nodeP->rplimit = Bcast_Cset (nodeP, xload, nodeP->threshold, 1);
   if (!nodeP->rplimit)
     nodeP->rplimit = bcast nbor(nodeP,
                  probe limit,
                  xload);
   free ((char *)dblock);
   /* Add low T KCALL time out message */
   nodeP->pwait_set = nodeP->rplimit;
   nodeP->low t set = 1;
   restart_clock(nodeP, LOW_TOUT,
           2 * probe limit * ARR_TIME + TIMEOUT,
           NULL):
  return TRUE;
```

```
ADAPTIVE SEND
                                                          */
         ***************
int Rev_avg_snd(PROCESSOR *nodeP, int xload)
   if (nodeP->high\ t\ set == 0)
    nodeP->high_t_set = 1;
    restart_clock(nodeP, HIGH_TOUT,
         2 * probe_limit * ARR_TIME + TIMEOUT,
   return TRUE:
}
       ***************************
    GLOBAL AVERAGE
int global_avg_snd(PROCESSOR *nodeP,int xload)
  if (nodeP->high_t_set == 0)
   dblock = pack_data (NULL, HIGH_LD_MSG, 4);
   if (Bcast_Cset(nodeP, xload, nodeP->threshold, 0));
    bcast_nbor(nodeP, probe limit, xload);
   free ((char *)dblock);
   /* Add high T KCALL time out message */
   nodeP->high_t_set = 1;
   restart_clock(nodeP, HIGH_TOUT,
         2 * probe_limit * ARR_TIME + TIMEOUT,
  return TRUE;
}
       GLOBAL_AVERAGE
   ***************
int global avg rcv(PROCESSOR *nodeP)
   if (nodeP->low\ t\ set == 0)
   nodeP->low_t_set = 1;
   restart_clock(nodeP, LOW_TOUT,
         2 * probe limit * ARR_TIME + TIMEOUT,
         NULL);
  return TRUE:
```

```
ADD_PROCESS
                                                               */
USER_PROCESS *Add_Process (int process_id,
            USER_PROCESS *proc P,
            PROCESSOR *node P)
 USER PROCESS
                      *pblock,
                      *create_new_process(int, PROCESSOR *);
 node_P->n_local_procs++;
 node_P->n_active_local_procs++;
 if (proc_P == NULL)
  pblock
          = create_new_process (process_id, node_P);
 else
   pblock
                = proc P;
   pblock->exec_here_time = -1;
  };
 if (pblock != NULL)
   pblock->itype
                  = 0:
   pblock->timeout
                   = 0;
   pblock->nxt_proc_ptr = NULL;
   if (node_P->firstjb_ptr == NULL)
    node_P->firstjb_ptr = pblock;
    node P->lastjb ptr = pblock;
   else
      node_P->lastjb_ptr->nxt_proc_ptr = pblock;
      node P->lastjb ptr
                           = pblock;
  return (pblock);
KREATE_FRAME
MSG_FRAME *Kreate_Frame(PROCESSOR *node_P,
            USER_PROCESS *msg,
            int sno)
{MSG_FRAME
                   *buffer;
         = msg_packet (node_P, node_P->node_id);
  buffer
                 = KC_CALL;
  buffer->type
                   = sno;
  buffer->service_no
                   = msg;
  buffer->mess_data
                   = 0;
  buffer->total bytes
  return buffer;
```

```
SCHEDULE_PROCESS
void schedule_process (PROCESSOR *node_P)
{ USER_PROCESS
                      *proc P,
              *curr process,
              *find_process(PROCESSOR *, USER_PROCESS *);
 int
               reverse_policy(PROCESSOR *, int);
 void
               kernal_msg_handler(PROCESSOR *, MSG_FRAME *);
 curr_process = node_P->curr_process;
 if ((curr_process != NULL)
   && (curr_process->exec_time >= EXECUTION_TIME))
    /* set up parameters */
    Del++;
    kernal_msg_handler (node_P,
              Kreate_Frame(node_P, curr_process, HALT_CALL));
    curr_process = Remove_Process(node_P, curr_process->upid);
    proc_P = node_P->curr_process;
    curr_process->nxt_proc_ptr = NULL;
    free ((char *) curr process);
    receiver initiated (node P);
   proc_P = find_process (node P, curr process);
  node P->curr process = proc P;
  if (proc P!= NULL)
    kernal_msg_handler (node_P,
              Kreate_Frame(node_P, proc_P, EXEC_CALL));
  }
  else
    time_update (node_P, (int) ELAPSED_TIME, OStime);
```

```
FIND_PROCESS
                  *find_process(PROCESSOR *node_P,
USER_PROCESS
                USER_PROCESS *proc_P)
 USER_PROCESS
                       *curr_P,
                       *prev_P;
 prev_P = curr_P = proc_P;
 do
  if (curr_P == NULL)
    curr_P = node_P->firstjb_ptr;
  else
     curr_P = curr_P->nxt_proc_ptr;
  if (curr_P == prev_P)
    break;
 }while ((curr P == NULL)
      || (curr_P->schedulable == FALSE)
      || (curr_P->blocked == TRUE));
 return curr_P;
}
```

```
/****************
/*
    REMOVE_PROCESS
                                                                     */
*Remove_Process ( PROCESSOR *node_P, int process_id)
 USER_PROCESS
                      *curr process,
                      *prev Pess;
 curr_process = node_P->firstjb_ptr;
 if (curr_process != NULL)
    prev_Pess = curr process;
    while (curr process != NULL)
     if (curr_process->upid == process id)
     else
      {
       prev_Pess = curr_process;
       curr_process = curr_process->nxt_proc_ptr;
    if (curr_process != NULL)
      if (curr_process == prev_Pess)
        node P->firstjb ptr = curr process->nxt proc ptr;
        prev_Pess = NULL;
      else
        prev_Pess->nxt_proc_ptr
          = curr_process->nxt_proc_ptr;
      if (node P->lastjb ptr == curr process)
        node P->lastjb_ptr = prev_Pess;
 if (curr_process != NULL)
   if (node P->curr_process == curr_process)
     node_P->curr_process = curr_process->nxt_proc_ptr;
     if ((node_P->curr_process != NULL)
       &&(node P->curr_process->nxt_proc_ptr == curr_process))
       node P->curr_process = curr_process->nxt_proc_ptr;
  return curr_process;
```

```
CREATE_NEW_PROCESS
/*
USER_PROCESS *create_new_process(int process_id,PROCESSOR *node_P)
 USER_PROCESS
                  *pblock;
         = (USER_PROCESS *)
 pblock
         malloc (sizeof(USER_PROCESS));
 if (pblock != NULL)
    pblock->upid
                  = process id;
    pblock->orig_mc
                   = node_P->node id;
   pblock->exec time
                    = 0;
   pblock->exec_here_time = 0;
   pblock->exist_time = 0;
   pblock->nxt_proc_ptr = NULL;
   pblock->schedulable = TRUE;
   pblock->blocked
                   = FALSE;
                   = FALSE;
   pblock->finished
   pblock->n_probes
 else
    fprintf (monitor_file,
     "\n\n::M A L L O C ****** FAILED");
    fprintf (monitor file,
          "in create_new_process()\n");
    fprintf (monitor_file, "Process No.%d\n", process_count);
    exit (0);
  }
 return pblock;
KERNAL_CALL
                                                            */
void kernal_call (PROCESSOR *node P, MSG FRAME *msg)
void
            do_processing(PROCESSOR *);
USER PROCESS
                  *create_new_process(int, PROCESSOR *);
switch (msg->service no)
 /* case CREATE_CALL: create_new_process();
         break; */
 case HALT CALL: exit proc(node P, msg);
         break;
 case EXEC CALL: do_processing(node_P);
         break;
 case RESET_TMER:
         break;
 case CANCEL TMER:
         break;
}}
```

```
E\:X\:T\:E\:R\:N\:A\:L\_\:M\:E\:S\:S\:A\:G\:E
                                                                       */
void external_message (PROCESSOR *node_P, MSG_FRAME *msg)
 switch (msg->service_no)
 {case LOAD_STATE:
       switch (lba index)
      {case 2:
       case 4:
       case 5: Update_Cset (node_P, msg, ASCEND);
            break;
        case 3: Update_Cset (node_P, msg, DESCEND);
       }
      break;
 case PSTOP_MSG: exit_msg(node_P, msg);
            break;
 case HIGH_LD_MSG: threshold = node_P->threshold;
 case PRB_MSG: if (probe_msg (node_P, msg))
             Update_Cset(node_P, msg, DESCEND);
           node_P->low_t_set = 0;
           break;
 case PRB_RPLY_MSG:
      Update_Cset (node_P, msg, ASCEND);
      switch (lba_index)
       case 2: probe_reply_msg (node_P, msg);
           break;
       case 4:
       case 5: lcal_reply_msg (node_P, msg);
           break;
       case 6:
       case 9:Gbal_reply_msg (node_P, msg);
      break;
case WK RPLY MSG:
     Update_Cset (node_P, msg, DESCEND);
     switch (lba_index)
       case 3:if (node_P->pwait_set)
            work_reply_msg (node_P, msg);
          break;
       case 7:
       case 8:
       case 9:
       case 10:Gbal_reply_msg (node_P, msg);
     break;
case PRC_ARI_MSG: process_msg (node_P, msg);
           break;
```

```
case LOW_LD_MSG:
  case WK_RQST_MSG: if (work_request_msg (node_P, msg))
              Update_Cset (node_P, msg, ASCEND);
            if (node_P->high_t_set)
              Cancel_Timer(HIGH_TOUT, node_P, NULL);
            break;
  case SND_PRC_MSG: transfer_msg (node_P, msg);
            break;
  case NEW_LD_MSG: new_load_msg (node_P, msg);
            break;
}
     DO_PROCESSING
<del>/***********************</del>
void do_processing (PROCESSOR *node_P)
 context switch;
 time_update(node P,
       (int) (CONTEXT_SWITCH * AVE_INST + TIME_SLICE * AVE_INST).
       USERtime);
}
/*
     HI_TIMEOUT_MSG
                                                                     */
/*
/************************
void hi_timeout_msg(PROCESSOR *nodeP)
 int
          xload;
 xload = nodeP->n_local_procs + nodeP->n_virtual_procs
     - nodeP->threshold;
 if (nodeP->high_t_set \&\& (xload > 0))
   nodeP->threshold++;
   if (lba_index >7)
    xload = nodeP->n_local_procs + nodeP->n_virtual procs;
           = pack_data (NULL, NEW LD MSG, 4);
    dblock -> Idv = (double) nodeP -> threshold;
    if (Bcast_Cset(nodeP, xload, nodeP->threshold - 1, 0));
      bcast_nbor(nodeP, probe_limit, xload);
    free ((char *)dblock);
 nodeP->high t set = 0;
#ifdef DEBUG MSG
 if ((process_count >= MIN_PID)
     && (process count <= MAX_PID))
 fprintf (monitor file,
      "\n HIGHLD TIMED OUT for Node%d at %8.3f\tNewTh=%3d\n",
```

```
nodeP->node_id,
      nodeP->sys_real_time/ONE_SECOND,
      (int) nodeP->threshold);
#endif
}
     LOW_TIMEOUT_MSG
<del>/***********************</del>
void low_timeout_msg(PROCESSOR *nodeP)
         xload;
 xload = nodeP->n_local_procs + nodeP->n_virtual_procs
       - nodeP->threshold;
  if (nodeP->low_t\_set \&\& (xload < 0))
  if (nodeP->threshold > 1)
    nodeP->threshold --;
    if (lba index > 7)
     xload = nodeP->n_local_procs + nodeP->n_virtual_procs;
     dblock = pack_data (NULL, NEW LD MSG, 4);
     dblock->ldv = (double)nodeP->threshold;
     if (Bcast_Cset(nodeP, xload, nodeP->threshold - 1, 1));
       bcast_nbor(nodeP, probe_limit, xload);
     free ((char *)dblock);
  }
 nodeP -> low_t_set = 0;
#ifdef DEBUG MSG
 if ((process_count >= MIN_PID)
     && (process_count <= MAX_PID))
 fprintf (monitor_file,
      "\n LOWLD TIMED OUT for Node%d at %8.3f\t NewTh=%3d\n",
      nodeP->node id,
      nodeP->sys real time/ONE_SECOND,
      (int)nodeP->threshold);
#endif
}
```

```
AVGLD_TOUT_MSG
                                                                      */
 void Avgld_tout_msg(PROCESSOR *nodeP)
  int
         xload;
  xload = nodeP->n_local_procs + nodeP->n_virtual_procs
       - nodeP->threshold;
  if (x load != 0)
    dblock = pack_data (NULL, NEW_LD_MSG, 4);
     nodeP->threshold = Cset_Load_Avg(nodeP);
    dblock->ldv = (double) nodeP->threshold;
     bcast_nbor(nodeP, probe_limit, nodeP->n_local_procs);
     free ((char *)dblock);
  }
  nodeP->low_t set = 0;
  nodeP->high t set = 0;
#ifdef DEBUG MSG
  if ((process_count >= MIN PID)
     && (process_count <= MAX_PID))
  fprintf (monitor_file,
      "\n RCV AVGLD TIMED OUT for Node%d at \%8.3 f\t NewTh=\%3 d\n",
      nodeP->node id,
      nodeP->sys_real_time/ONE SECOND,
      (int)nodeP->threshold);
#endif
}
     WRQ_TIMOUT_MSG
void wrq_timout_msg (PROCESSOR *node ptr)
#ifdef DEBUG_MSG
 if ((process count >= MIN PID)
     && (process_count <= MAX_PID))
 fprintf (monitor file,
      "\n WAITP TIMED OUT for Node%d at %8.3f\n",
     node_ptr->node_id,
     node_ptr->sys_real_time/ONE_SECOND);
#endif
 /* node ptr->pwait set++;
 node ptr->n virtual_procs--;
}
```

```
PRB_TIMOUT_MSG
void prb_timout_msg (PROCESSOR *node_ptr, USER_PROCESS *prc)
 /* NO REPLY RECEIVED */
 node_ptr->rplimit--;
#ifdef DEBUG_MSG
 if ((process_count >= MIN_PID)
    && (process_count <= MAX_PID))
 printf ("\n PROBE TIMED OUT for prc%d \n", prc->upid);
#endif
 node_ptr->rplimit--;
 send_probe(node_ptr, prc);
/*****************
    SRQ_TIMOUT_MSG
/**********************
void srq_timout_msg (PROCESSOR *node_ptr,USER_PROCESS *prc)
/* -- unblock migratable process
 -- reset node_P->SRQ_TIMOUT
#ifdef DEBUG_MSG
 if ((process_count >= MIN_PID)
    && (process_count <= MAX_PID))
 fprintf (monitor file,
     "\n SRQ TIMED OUT for Node%d at %8.3f\n",
    node ptr->node id,
     node_ptr->sys_real_time/ONE_SECOND);
#endif
 prc->schedulable = TRUE;
 prc->itype
            = 0;
 prc->timeout = 0;
 node_ptr->n_active_local_procs++;
 if (node ptr->rplimit > 0)
   node ptr->rplimit--;
}
```

```
CANCEL_TIMER
BOOLEAN Cancel_Timer (int itype,PROCESSOR *nodeP,USER_PROCESS *prc)
 USER_PROCESS *this_proc;
 if (prc == NULL)
  if (itype == nodeP->timq.itype)
   /* re-initialise */
    switch (nodeP->timq.itype)
     case WAITP_TOUT: nodeP->rplimit = 0;
              nodeP->pwait set = 0;
              nodeP->n_virtual_procs = 0;
     case HIGH\_TOUT : nodeP->high\_t\_set = 0;
              break;
     case LOW TOUT:
             nodeP->low_t set = 0;
             break;
    case RCV_AVG_TOUT:
             nodeP->high_t_set = 0;
             nodeP->low_t_set = 0;
             break;
   Select_next timeout(nodeP);
   return TRUE;
 else
  this proc
              = nodeP->firstjb ptr;
  while (this_proc != NULL)
   if (prc == this_proc)
     if (prc->itype == itype)
      /* re-initialise */
      prc->timeout = 0;
      prc > itype = 0;
      if (nodeP->timq.minP == prc)
      Select next timeout(nodeP);
      return TRUE;
     break;
   else
    this proc = this proc->nxt_proc_ptr;
return FALSE;
```

```
#include <stdio.h>
#include <stdlib.h>
#include <signal.h>
#include <unistd.h>
#include <strings.h>
#include "model_params.h"
#include "model_types.h"
#include "globvar.h"
#define SYNTH_WORKLOAD
                            2
GLBAL FRAME
                   *dblock:
MSG_FRAME
                  *create_frame(PROCESSOR *, int);
ROUTE_ENTRY route_node (short int, LINK_ENTRY [], ROUTE_ENTRY []);
extern int
           SkipRun;
extern NODE_ENTRY
                   *processor_table;
extern int
             probe limit,
             process count,
             total_pe;
extern int
             FinL, MaxL;
extern float
              threshold;
extern FILE
              *monitor_file;
TRANSMIT
                                                              */
  void transmit (PROCESSOR *sender,
       ROUTE_ENTRY nxt_node,
       MSG_FRAME *msg_P)
 sender->n tx++;
 msg P->time_stamp += PROTOCOL_TIME * AVE_INST
      + msg_P->total_bytes * TX_BYTE_TIME;
 Write_To_MsgQ(&(nxt_node.node_ptr->queue), msg_P);
 time_update (sender, (int) (PROTOCOL_TIME * AVE_INST
       + msg_P->total_bytes * TX_BYTE_TIME), OStime);
```

```
RECEIVE
/****************
MSG_FRAME *receive ( PROCESSOR *node_P)
 ROUTE ENTRY
                      next route;
 MSG FRAME
                      *msg P;
 double
               most_recent_time;
 most_recent_time = node_P->sys_real_time;
 while ((node_P->queue.head != NULL) &&
               (most_recent_time >= node_P->queue.head->time_stamp))
       msg_P = Read_From_MsgQ(&(node_P->queue));
       node_P->n_rx++;
        time_update (node_P, (int) (PROTOCOL_TIME * AVE_INST
               + msg_P->total_bytes * RX_BYTE_TIME), OStime);
        if (msg_P->dest_id == node P->node id)
#ifdef DEBUG MSG
  if ((process_count >= MIN_PID)
     && (process_count <= MAX_PID))
    fprintf(monitor_file, "\nNode%d receives Packet N%d-N%d::\t",
       node_P->node_id, msg_P->sender_id, msg_P->dest id);
#endif
               return (msg_P);
     }
        else
         msg_P->distance++;
         next_route = route_node (msg_P->dest_id,
                    node_P->links,
                    node P->route table);
         transmit (node P, next route, msg P);
 }/* end while */
 return (NULL);
}
```

```
CREATE_FRAME
                                                                  */
MSG_FRAME *create_frame(PROCESSOR *node_P, int dest_id)
 MSG_FRAME
                  *buffer;
   if ((buffer = msg_packet (node_P, dest_id)) != NULL)
    buffer->type
                   = dblock->type;
    buffer->service_no = dblock->sno;
    buffer->mess_data = dblock->prc;
    buffer->ldvalue
                  = dblock->ldv;
    buffer->total_bytes = dblock->size;
   return buffer;
 }
    msg_packet
MSG_FRAME *msg_packet (PROCESSOR *node_P, int dest_id)
 MSG_FRAME *buffer;
        = (MSG FRAME *)
          malloc (sizeof(MSG_FRAME));
 if (buffer != NULL)
  buffer->dest_id
                   = dest id;
  buffer->sender_id = node_P->node_id;
  buffer->time_stamp = node_P->sys_real_time;
  buffer->distance
  buffer->total bytes = 0;
  buffer->next_pack_ptr = NULL;
  return (buffer);
 }
else
  { /* ** ANSI Restriction ** *
    longjmp (ErrLabel, 3);
   ** **
             */
 }
}
```

```
pack_data
 GLBAL_FRAME *pack_data (USER_PROCESS *proc, int service, int sz)
 GLBAL FRAME
                 *buffer:
 buffer
         = (GLBAL_FRAME *)
           malloc (sizeof(GLBAL_FRAME));
 if (buffer != NULL)
   buffer->type
                    = KC EMSG;
   buffer->sno
                   = service;
   buffer->prc
                   = proc;
   buffer->size
                   = sz;
 }
 else
  {
     fprintf (monitor file,
      "\n\n::M A L L O C ****** FAILED");
     fprintf (monitor_file,
            "in pack_data()\n");
     fprintf (monitor_file, "Process No.%d\n", process_count);
     exit (2);
  }
 return (buffer);
}
/*
    MIGRATABLE_TASK
/*****************
USER_PROCESS
                *migratable_task( PROCESSOR *nodeP)
 USER_PROCESS
                      *proc;
 BOOLEAN
                   unsuitable;
        = nodeP->firstjb_ptr;
 unsuitable = (proc == NULL) || (proc->schedulable == FALSE)
        || (proc->blocked == TRUE)
        || (proc == nodeP->curr_process)
        || (proc->exec_here_time < 0);</pre>
 do
  if ((proc == NULL) || (!unsuitable))
   break;
  proc = proc->nxt_proc_ptr;
  unsuitable = (proc == NULL) || (proc->schedulable == FALSE)
       || (proc->blocked == TRUE)
       || (proc == nodeP->curr_process)
       \parallel (proc->exec here time < 0);
 } while (unsuitable);
return proc;
```

```
LONGEST_RUNNER
                                                                         */
USER_PROCESS
                 *longest_runner(PROCESSOR *nodeP)
 USER PROCESS
                       *proc, *lrp;
 BOOLEAN
                    unsuitable;
         = nodeP->firstjb_ptr;
 proc
 lrp
        = NULL;
 unsuitable = (proc == NULL) \parallel (proc->schedulable == FALSE)
        || (proc->blocked == TRUE)
        || (proc == nodeP->curr_process)
        || (proc->exec_here_time < 0);</pre>
 do
  if (proc == NULL)
    break;
  if (!unsuitable)
       if ((lrp == NULL)
      \parallel (proc\text{->}exec\_time < lrp\text{->}exec\_time))
     Irp = proc;
  proc = proc->nxt_proc_ptr;
 unsuitable = (proc == NULL) || (proc->schedulable == FALSE)
        || (proc->blocked == TRUE)
        || (proc == nodeP->curr_process)
        || (proc->exec_here_time < 0);
 } while (1);
return lrp;
```

```
ROUTE_SELECTION
                                                                       */
ROUTE_ENTRY
                  Route_Selection(int dn,
                 LINK_ENTRY links[],
                 ROUTE_ENTRY rtable[])
 int
        nbor_id, pe,
       km, ridx, i;
 nbor_id = links[dn].node id;
     = links[dn].hops;
 for (i = 0; i < TABLE\_SIZE; i++)
  pe = rtable[i].node_id;
  if (pe == EMPTY)
    continue;
   if (pe == nbor id)
     ridx = i;
   else
   if (processor_table[pe].node_adr->links[dn].hops
      == km - 1)
     ridx = i;
     break;
  links[dn].node_id = rtable[ridx].node id;
  return rtable[ridx];
}
    ROUTE_NODE
                                                                      */
ROUTE_ENTRY route node (short int dest node,
             LINK_ENTRY links[],
             ROUTE_ENTRY rtable[])
int
       idx;
LINK_ENTRY actual_cell;
actual_cell = links[dest_node];
for (idx = 0; idx < TABLE\_SIZE; idx++)
   if (rtable[idx].node_id != EMPTY)
    if (actual cell.node id == rtable[idx].node_id)
      break;
return rtable[idx];
```

```
FIND_NBOUR
                                                                    */
 int find_nbour (PROCESSOR
                            *pe, int maxd)
  static int
              i = 0:
  auto int
              j;
  if (i == total_pe)
   i = 0;
  j = i;
  do
    if (pe->links[i].hops == maxd)
     j = i;
     i++;
     return j;
    else i = (i + 1) \% total pe;
  } while (j != i);
  return -1;
 /*
/*
     MIGRATE
/*
void migrate(PROCESSOR *node P,
         USER_PROCESS *proc_P.
        int dest id)
{ ROUTE ENTRY
                       next_route;
   proc P->itype
                   = 0;
   proc_P->timeout
                    = 0;
   next_route
                  = route_node ((short)dest_id,
                        node P->links,
                        node_P->route_table);
  if (proc_P == NULL) return;
  time_update (node_P, (int) TIME_MIG_PROC, OStime);
#ifdef DEBUG MSG
  if ((process_count >= MIN_PID)
     && (process count <= MAX PID))
    fprintf (monitor file,
        "\nnode %d sends process%d [load = %d] to node %d\n",
       node_P->node id,
       proc P->upid,
       node P->load, dest id);
#endif
  dblock = pack_data (proc_P, PRC_ARI_MSG, 10218);
  node_P->n_migrates++;
  transmit (node_P,
       next route,
       create frame (node P, dest id));
  free ((char *)dblock);
```

```
PROBED
BOOLEAN probed (USER_PROCESS *proc_P, int mc_id)
BOOLEAN found;
found = FAILED;
if (proc_P->n_probes > 0)
 int idx;
 idx = 0;
 while (idx <= probe_limit -1)
 if ((proc_P->mcs_probed)[idx] == mc_id)
  found = SUCCESS;
  break;
 idx++;
};
return found;
```

```
SEND_PROBE
void send_probe (PROCESSOR *node_P, USER_PROCESS *proc_P)
 if (proc_P->n_probes < probe_limit)
 {
  int mc_id,
     idx;
  idx = proc_P->n probes;
  if ((node_P->cSet[0].id < 0)
    \parallel (node_P->cSet[0].ld > 0))
   do
   {
    mc_id = randmc_id(node_P);
   } while ((probed(proc_P, mc_id))
        || (mc_id == node_P->node_id));
  else mc_id = node_P->cSet[0].id;
  proc_P->mcs_probed[idx] = mc_id;
  proc_P->n_probes++;
  /* send probe message */
  int dstance, load;
  MSG_FRAME *buff;
  dblock = pack_data (proc_P, PRB_MSG, 4);
  buff = create_frame (node_P, mc_id);
  load = (int)(node_P->n_local_procs - N_SYS_PROCS - threshold);
  buff->ack_reply = load;
  transmit (node_P,
        route_node (mc_id, node_P->links,
                 node_P->route_table),
       create_frame (node_P, mc_id));
  free ((char *)dblock);
  dstance = node_P->links[mc_id].hops;
  restart_clock (node_P, PRB_TOUT, (double)
                 ARR_TIME * dstance * 2 + TIMEOUT,
          proc P);
  node P->rplimit++;
else /* probe_limit has been exceeded */
  node_P->n_active_local_procs++;
 proc P->schedulable = TRUE;
```

```
BROADCAST
             void broadcast (PROCESSOR *node_P)
{
 int
               i;
 for (i = 0; i < total pe; i++)
  if (node_P->node_id == i)
   continue;
  else
   /* send message */
    transmit (node_P,
         route_node (i, node_P->links, node_P->route_table),
         create_frame (node P, i));
 }
}
BOOLEAN Bcast_Cset(PROCESSOR *pe, int load, int threshold, int thold)
  MSG FRAME *buff;
  int i = 0, send, tx = 0;
  while ((i < CSET\_SIZE) \&\& (pe->cSet[i].id > 0))
    if (thold == 0)
      send = pe->cSet[i].ld < threshold + N_SYS_PROCS;
    else
      send = pe->cSet[i].ld > threshold + N_SYS_PROCS;
    if (send)
     /* send probe */
     tx++;
     buff = create_frame (pe, pe->cSet[i].id);
     buff->ack_reply = load;
     transmit (pe,
           route_node (pe->cSet[i].id,
                  pe->links,
                  pe->route_table),
           buff);
    i++;
   return tx;
```

```
BCAST_NBOR
int bcast_nbor (PROCESSOR *pe, int max_d, int load)
 int
       i, d;
 int
       total = 0;
 MSG_FRAME *buff;
 for (i = 0; i < total_pe; i++)
  d = pe->links[i].hops;
  if ((d > 0) \&\& (d \le max_d))
   /* send probe message */
    total++;
    buff = create_frame (pe, i);
    buff->ack_reply = load;
    transmit (pe,
         route_node (i, pe->links, pe->route_table),
  };
 }
return total;
    NEW_LOAD_MSG
void new load msg (PROCESSOR *nodeP, MSG FRAME *msg)
 nodeP->low_t_set = 0;
 nodeP->high_t_set = 0;
 Cancel_Timer(LOW_TOUT, nodeP, NULL);
 Cancel_Timer(HIGH_TOUT, nodeP, NULL);
 if (msg->ldvalue > MaxL)
   MaxL = (int) msg->ldvalue;
 FinL = (int) msg->ldvalue;
 nodeP->threshold = (int) msg->ldvalue;
```

```
PROBE_MSG
                                                                         */
/*
      ****************************
BOOLEAN probe_msg (PROCESSOR *node_P, MSG_FRAME *msg)
 {
 MSG_FRAME
                      *buff;
               xload:
 ROUTE_ENTRY
                       next route:
 extern int
                 MsgQ_Length(),
              ProcQ_Length();
 /* received by CPU(i) */
  dblock = pack_data (msg->mess_data, PRB_RPLY_MSG, 5);
         = create_frame (node_P, msg->sender_id);
#ifdef DEBUG_MSG
  if ((process_count >= MIN_PID)
     && (process_count <= MAX_PID))
   fprintf (monitor file,
        "Node%d -overloaded- has asked node%d TO acceptn",
        msg->sender id,
        node P->node id);
   fprintf (monitor_file, "msqQL=%d procQL = %d\n",
       MsgQ_Length(node_P->queue.head),
       ProcQ Length(node P->firstjb_ptr));
   Dump_Msg(msg);
#endif
  xload = node_P->n_local_procs + node_P->n_virtual procs;
  buff->ack\_reply = xload;
  if (buff->ack reply <= threshold)
  { /* reply I am prepared to accept */
    buff->ack_reply = -(buff->ack_reply);
       node_P->n_virtual_procs++;
    node_P->rplimit++;
       restart_clock(node P, WAITP TOUT,
                  Delay_Time, NULL);
  }
  next_route
             = route node (msg->sender id,
                        node P->links,
                        node_P->route_table);
  transmit (node_P, next_route, buff);
  free ((char *)dblock);
  return xload <= threshold;
```

```
WORK_REQUEST_MSG
          ****************
BOOLEAN work_request_msg (PROCESSOR *node_P,MSG_FRAME *msg)
  MSG_FRAME
                       *buffer:
  USER_PROCESS
                        *prc,
               *migratable_task(PROCESSOR *);
                xload;
 ROUTE ENTRY
                        next_route;
 /* -- find a process that has not commenced execution
   -- find a process that is blocked
   -- or find a process to pre-empt
   -- Block it
 */
          = NULL;
  prc
           = pack_data (msg->mess_data, WK_RPLY_MSG, 5);
  buffer
           = create_frame (node_P, msg->sender_id);
           = node_P->n_local_procs + node_P->n_virtual_procs;
  xload
#ifdef DEBUG_MSG
  if ((process_count >= MIN PID)
     && (process_count <= MAX_PID))
   fprintf (monitor file,
        "Node%d has asked node%d for WORK::::",
        msg->sender_id, node_P->node_id);
   fprintf \ (monitor\_file, \ "\ \ \ \ \ \ \ \ procQL = \%d\ \ \ \ \ )
       MsgQ_Length(node_P->queue.head),
       ProcQ_Length(node_P->firstjb_ptr));
  }
#endif
  buffer->ack_reply = xload;
  if (xload > node_P->threshold + N_SYS_PROCS)
    prc = longest_runner (node_P);
         if (prc != NULL)
     { /* reply - yes I have processes I wish to migrate
       buffer->mess data = prc;
      prc->schedulable = FALSE;
      node_P->n active local procs--;
      restart_clock(node_P, SRQ_TOUT, Delay Time, prc);
      node P->rplimit++;
      */
     } else buffer->ack_reply = -(buffer->ack_reply);
  } else buffer->ack_reply = -(buffer->ack_reply);
  next_route = route node (msg->sender id,
                 node P->links,
                 node P->route_table);
  transmit (node P, next route, buffer);
  free ((char *)dblock);
  return xload > node P->threshold + N_SYS PROCS;
```

*/

```
PROBE_REPLY_MSG
                                                                */
 /*
 void probe_reply_msg (PROCESSOR *node_P, MSG_FRAME *msg_P)
  USER_PROCESS
                     *proc_P;
  /* received by CPU(i) */
 #ifdef DEBUG_MSG
   if ((process_count >= MIN PID)
      && (process_count <= MAX_PID))
    fprintf (monitor_file,
        "Node%d::Time:%f\tNode%d's reply (prc%d) was %d\n",
        node_P->node_id,
        node_P->sys_real_time / ONE_SECOND,
        msg_P->sender_id, msg_P->mess_data->upid,
        msg_P->ack_reply);
 #endif
 if \ (Cancel\_Timer(PRB\_TOUT,node\_P, \ msg\_P->mess\_data)) \\
    node_P->rplimit--;
    if (msg_P->ack_reply > FALSE)
     send_probe (node_P, msg_P->mess_data);
    else
     int idx;
     proc_P = msg_P->mess_data;
     proc_P = Remove_Process (node_P, proc_P->upid);
     for (idx = 0; idx < proc_P -> n_probes; idx++)
      proc_P->mcs_probed[idx] = EMPTY;
     proc_P->n_probes = 0;
     migrate (node_P, proc_P, msg_P->sender_id);
    node_P->n_local_procs--;
  }
}
```

```
WORK_REPLY_MSG
     void work_reply_msg (PROCESSOR *node_P, MSG_FRAME *msg)
 MSG FRAME
                     *buffer;
 ROUTE_ENTRY
                      next_route;
              xload;
#ifdef DEBUG MSG
  if ((process_count >= MIN_PID)
     && (process_count <= MAX_PID))
   fprintf (monitor file,
        "Node%d::Time:%f\tNode%d's reply was %d\n",
        node_P->node_id,
        node_P->sys_real_time / ONE SECOND,
        msg->sender id,
        msg->ack_reply);
#endif
  if (msg->ack_reply > FALSE)
  {
   dblock = pack_data (msg->mess_data, SND_PRC_MSG, 5);
   buffer = create_frame (node P, msg->sender id);
   xload = node_P->n_local_procs + node_P->n_virtual_procs;
   buffer->ack\_reply = xload;
   if (xload <= threshold)
    buffer->ack_reply = -(buffer->ack_reply);
    node_P->n_virtual_procs++;
/*
        if (node_P->timq.mintime > node_P->sys_real_time + Delay_Time)
         node P->timq.mintime = node P->sys real time + Delay Time;
*/
        restart_clock(node_P, WAITP_TOUT,
                        Delay_Time,
              NULL);
   node P->pwait set--;
               = route node (msg->sender_id,
   next_route
                   node_P->links,
                   node P->route_table);
   transmit (node_P, next_route, buffer);
}
```

```
LCAL_REPLY_MSG
void lcal_reply_msg ( PROCESSOR *node_P, MSG_FRAME *msg_P)
USER_PROCESS
                   *proc P;
/* received by CPU(i) */
if (node_P->rplimit)
 node P->rplimit--;
 if (msg_P->ack_reply <= FALSE)
   proc_P = msg_P->mess_data;
  /* Cancel Timer */
  if (Cancel_Timer(SRQ_TOUT, node_P, proc_P))
    node_P->rplimit = 0;
    proc_P = Remove_Process (node_P, proc_P->upid);
    if (proc_P != NULL)
    migrate (node_P, proc_P, msg_P->sender_id);
    node_P->n_local_procs--;
```

```
*/
     GBAL_REPLY_MSG
                                                                     */
 /*********************
 void Gbal_reply_msg (PROCESSOR *node_P, MSG_FRAME *msg_P)
  USER_PROCESS
                      *proc_P;
 /* received by CPU(i) */
 switch (msg_P->service no)
  {
  case PRB_RPLY_MSG:
     if (node_P->high_t_set &&
       (msg_P->ack_reply < FALSE))
      Cancel\_Timer(HIGH\_TOUT, node\_P, NULL);
      Cancel_Timer(RCV_AVG_TOUT, node_P, NULL);
      node_P->high_t_set=0;
       proc_P = longest_runner (node_P);
      if (proc_P != NULL)
       migrate (node_P,
            Remove_Process (node_P, proc_P->upid),
            msg_P->sender_id);
       node_P->n_local_procs--;
      break;
 case WK_RPLY_MSG:
     if (node_P->low_t_set
       && (msg_P->ack_reply > FALSE))
       Cancel\_Timer(LOW\_TOUT, node\_P, NULL);
       /* Cancel_Timer(RCV_AVG_TOUT, node_P, NULL); */
       node_P->timq.itype = WAITP_TOUT:
       threshold = node_P->threshold;
       work_reply_msg (node_P, msg_P);
       node_P->low_t_set=0;
 }
}
```

```
TRANSFER_MSG
void transfer_msg (PROCESSOR *node_P, struct comm_packet *msg_P)
USER_PROCESS *proc_P;
/* received by CPU(i)
 if (node_P->rplimit)
  node P->rplimit--;
  proc_P = msg_P->mess data;
  Cancel_Timer(SRQ_TOUT, node_P, proc_P);
  if (msg_P->ack_reply < TRUE)
    proc_P = Remove_Process (node_P, proc_P->upid);
    if (proc_P != NULL)
     migrate (node_P, proc_P, msg_P->sender_id);
     node_P->n_local_procs--;
   else node_P->n_active_local_procs++;
   Select_next_timeout(node_P);
 }
 else
  if ((msg_P->ack_reply < TRUE)
     && (node_P->n_local_procs
       > node_P->threshold + N_SYS_PROCS))
      proc_P = longest_runner(node_P);
      if (proc_P != NULL)
       migrate (node P,
           Remove_Process (node_P, proc_P->upid),
           msg_P->sender_id);
       node_P->n_local_procs--;
}
```

```
PROCESS_MSG
 void process_msg (PROCESSOR *node_P,MSG_FRAME *msg)
  USER_PROCESS
                      *proc_P;
  proc_P = msg->mess_data;
  proc_P->exec_here_time = 0;
  proc_P->residency time = 0;
  proc_P->exist_time += msg->distance
                * (PROTOCOL_TIME * AVE_INST
                + msg->total_bytes * TX_BYTE_TIME);
  proc_P->schedulable = TRUE;
  proc_P->nxt_proc_ptr = NULL;
  (void) Add_Process (proc_P->upid, proc_P, node_P);
  node_P->n_immigs++;
 time_update (node_P, (int) TIME_RCV_PROC, OStime);
 #ifdef DEBUG MSG
   if ((process_count >= MIN_PID)
      && (process_count <= MAX_PID))
    fprintf (monitor_file,
        "\nnode %d RECEIVES process%d [load = %d] FROM node %d\n",
        node P->node id,
        proc_P->upid,
        node P->load, msg->sender_id);
#endif
 if (node_P->n_virtual_procs > 0)
  node_P->n_virtual_procs--;
  if (node_P->n_virtual_procs <= 0)</pre>
    Cancel_Timer(WAITP_TOUT, node_P, NULL);
 };
}
    EXIT_MSG
/********
void exit_msg (PROCESSOR *nodeP, MSG_FRAME *msg)
/* Do nothing */
nodeP->n deaths++;
nodeP->cumul_exist_time += msg->ldvalue;
}
```

```
EXIT PROC
                                                                    */
void exit_proc ( PROCESSOR *node_P, MSG_FRAME *msg)
{
 ROUTE ENTRY
                      next route;
 USER PROCESS
                     *proc P;
 extern void
                 exit_dump(PROCESSOR *, USER_PROCESS *);
 context_switch;
 time_update(node P,
       (int) (CONTEXT_SWITCH * AVE_INST + TIME_EXIT * AVE_INST),
       USERtime);
 proc_P = msg->mess_data;
 if (proc_P->orig_mc != node_P->node id)
  /* set up exit msg for original mc */
  /* send msg */
  /* TX (sizeof(EM_EXIT, OStime); */
  next_route = route_node (proc P->orig mc,
                node P->links,
                node_P->route table);
  dblock
           = pack_data (proc_P, PSTOP MSG, 4);
  dblock->ldv = proc P->exist time;
  transmit (node P.
       next route,
       create_frame (node_P, proc_P->orig_mc));
  free ((char *)dblock);
else
 node_P->cumul_exist_time += proc_P->exist_time;
 node P->n deaths++;
exit_dump(node P, proc P);
/* MARK process for removal */
proc_P->finished = TRUE;
proc P->schedulable = FALSE;
node_P->n_local_procs--;
node_P->n_active_local_procs--;
```

```
This is the the computational module, encompassing the accumulation
  of basic statistics such as averages and totals.
  It differs from the original by encompassing "ave_compute_or_display(),
  and the "job_creation() modules also.
  The job_creation module now provides a different randomizing seed
  each time the job is run. AND,
  Convergence is represented purely in terms of the percentage diff.
  in runtime between the fastest and slowest processor.
        Modified:
                       5th December 1991
 #include <stdio.h>
#include <stdlib.h>
#include <math.h>
#include "model params.h"
#include "model_types.h"
#include "globvar.h"
int
            Least loaded,
             last t,
             Most_loaded,
             sigma diff;
float
            Mean load,
             sigma_load,
             sigma var,
             sigma rtime,
             sigma_migrate,
             sigma_Tx,
             Mprofile,
             Ave migration(void),
             Ave_transmission(void);
extern float
            Estimator();
extern int
            njobs,
            probe limit,
             process count,
            total_pe;
extern float
            load value,
            threshold;
               lba index;
extern int
extern double
                 global_clock;
extern FILE
                *trace.
                   *seedfile,
            *monitor_file,
            *configfile.
            *paramfile:
                      *processor table;
extern NODE ENTRY
```

Sec. Del:

extern int

```
ave_compute_or_display
void ave_compute_or_display (BOOLEAN run complete)
 if (run_complete)
  float
          tsec;
  tsec = global_clock / ONE SECOND;
  fprintf(monitor_file, "\n%8.2f\t", tsec);
  fprintf(monitor_file, "%6.4f\t", sigma_var / tsec);
  fprintf(monitor_file, "%6.4f\t", sigma_load / tsec);
  fprintf(monitor_file, "%6.4f\t", sigma_diff / tsec);
  fprintf(monitor_file, "%6.4lf\t", Mprofile/tsec):
  fprintf(monitor_file, "%7.2f\t",
      100 * data_convergence((int) tsec));
  fprintf(monitor_file, "%7.4f\t", Ave_migration() / tsec);
  fprintf(monitor_file, "%7.4f\t", Ave_transmission()/ tsec);
  fprintf(monitor\_file, "\%5.2f\t\n", 100.0 * process\_count/ njobs);
 }
 else
  {
 /* fprintf (trace, "\t%2.2f\t%d\t%2.2f\t%2.2f\t%2.2f\t%2.2f\n",
  sigma_load / tsec,
  sigma_diff / tsec,
  sigma_var / tsec,
  Mprofile/tsec,
  sigma_migrate / tsec,
  sigma_Tx / tsec);
#ifdef DEBUG_MSG
  Dump_RTime();
#endif
}
```

```
*/
    final_compute
void final_compute (void)
 float
         tsec;
  tsec = global_clock / ONE_SECOND;
 fprintf(monitor\_file, "\n\%8.2f\t\%6.4f\t\%6.4f\t\%6.4f\t\%6.4f\t\%7.2f\t\%7.4f\t\%7.4f\t\%5.2f\t\n",
  sigma_var / tsec,
  sigma_load / tsec,
  sigma_diff / tsec,
  Mprofile/tsec,
  100 * data_convergence((int) tsec),
  Ave_migration() / tsec,
  Ave_transmission() / tsec,
  100.0 * process_count/ njobs);
#ifdef DEBUG MSG
  Dump_RTime();
#endif
 }
```

```
MEAN_DIFFERENCE
                                                                          */
float mean_difference (void)
 PROCESSOR
                     *nodeP:
 int
              total_load,
              idx;
 Least_loaded = 1000;
 Most loaded = 0;
 total load = 0;
 for (idx = 0; idx < total pe; idx++)
   fprintf (trace, "\n");
   nodeP = processor_table[idx].node_adr;
    fprintf (trace, "Node%d Load = %d ",
         nodeP->node_id, nodeP->n_local_procs);
*/
    if (nodeP->load < Least_loaded)</pre>
      Least_loaded = nodeP->load;
    if (nodeP->load > Most loaded)
      Most_loaded = nodeP->load;
    total load
                 += nodeP->load;
   fprintf (trace, "\n");
*/
return ((float)total load / total pe);
     VARIANCE
                                                                          */
/*************
float variance (float mean)
PROCESSOR
                     *nodeP;
float
               sum var,
              diff:
int
              idx;
            = 0.0;
sum_var
for (idx = 0; idx < total_pe; idx++)
   nodeP = processor_table[idx].node_adr;
             = ((float)nodeP->load) - mean;
   diff
   sum_var
               += (diff * diff);
return (sum_var / (total_pe - 1));
```

```
PROCESS_RUNTIME
float Process_runtime (void)
 PROCESSOR
                   *nodeP;
              sum_rtime,
 float
             rtime;
             idx;
 int
 sum_rtime
             = 0.0;
 for (idx = 0; idx < total_pe; idx++)
   nodeP = processor_table[idx].node_adr;
    if (nodeP->n_deaths != 0)
      rtime = nodeP->cumul_exist_time / nodeP->n_deaths;
      sum_rtime += (rtime / ONE_SECOND);
 return (sum_rtime / total_pe);
                                                                     */
     AVE_MIGRATION
float Ave_migration (void)
                    *nodeP;
  PROCESSOR
               sum migrate;
  float
              idx;
  int
                = 0.0;
  sum_migrate
  for (idx = 0; idx < total_pe; idx++)
    nodeP = processor_table[idx].node_adr;
    sum_migrate += nodeP->n_migrates;
  return (sum_migrate / total_pe);
```

```
*/
    AVE_TRANSMISSION
float Ave_transmission (void)
PROCESSOR
                 *nodeP;
float
            sum tx;
int
            idx;
         = 0.0;
sum tx
for (idx = 0; idx < total_pe; idx++)
  nodeP = processor_table[idx].node_adr;
  sum_tx += nodeP->n_tx;
return (sum_tx / total_pe);
        ************
    SIMULATION_COMPLETE
                                                              */
/*****************
BOOLEAN simulation_complete(void)
 BOOLEAN
             any_queue_empty();
 float
             data_convergence(int);
 int
               i, tsec;
tsec = (int)(global clock / ONE SECOND);
 i = tsec % SAMPLE_T;
 if (i!= last t)
 if (any_queue_empty())
   return QUEUES_FLUSHED;
  Mprofile += Process_runtime();
 Mean_load = mean_difference ();
 sigma_load += Mean_load;
 sigma_diff += (Most_loaded - Least_loaded);
 sigma var += variance(Mean_load);
 if ((i != last_t) && (tsec > STABLE_STATE))
  if ((tsec \% 10) == 0)
  if (data_convergence(tsec) < CONVERGE_FACTOR)
    return DATA_CONVERGED;
  else if ((Sec > 0) \&\& (process\_count/Sec > Del/Sec))
      return QUEUES_FLUSHED;
  }
last t = i;
 return FALSE;
```

```
*/
   data_convergence
data_convergence (int tsec)
float
        rtime,
             G_avg;
rtime = Process runtime();
G_avg = Mprofile/tsec;
if (rtime > G avg)
  return (rtime - G_avg) / rtime;
else
  return (G_avg - rtime) / G_avg;
*/
    TIME_UPDATE
void time_update (PROCESSOR *nodeP, int elapsed_time, int time_type)
 void
         quanta update(PROCESSOR *, int, int);
 float
         calc ave load(PROCESSOR *);
 USER_PROCESS *process_ptr;
 nodeP->sys_real_time += elapsed_time;
 if (time type == USERtime)
  nodeP->curr process->exec time
                              += elapsed time;
  nodeP->curr_process->exec_here_time += elapsed_time;
/*
 printf ("\n Time update P%d exec_time %lf\n",
    nodeP->curr process->upid,
    nodeP->curr_process->exec_time);
*/
 process ptr = nodeP->firstjb_ptr;
 while (process_ptr != NULL)
 process_ptr->exist_time += elapsed_time;
                = process_ptr->nxt_proc_ptr;
 process_ptr
 quanta_update (nodeP, elapsed_time, time_type);
 nodeP->load = (int) calc_ave_load (nodeP);
```

```
*/
    QUANTA_UPDATE
void quanta_update (PROCESSOR *nodeP, int elapsed time, int time_type)
 int added time,
   idx:
 void next_quantum(PROCESSOR *),
   add_to_quantum(PROCESSOR *, int, int);
 idx = nodeP -> quanta idx;
 while (elapsed_time > 0)
   added_time = elapsed time;
   if (elapsed_time + nodeP->quanta[idx].used >= QUANTUM)
    added_time = QUANTUM - nodeP->quanta[idx].used;
    add_to_quantum (nodeP, added_time, time_type);
    next_quantum(nodeP);
   else
     add_to_quantum (nodeP, added_time, time_type);
   elapsed_time -= added_time;
 }
}
ADD_TO_QUANTUM
/***************
void add_to_quantum (PROCESSOR *nodeP, int add_time, int time_type)
int idx;
idx = nodeP->quanta_idx;
 if (time type == OStime)
   nodeP->quanta[idx].OSportion += add_time;
   nodeP->OSoverhead
                       += add_time;
 nodeP->quanta[idx].used += add_time;
```

```
*/
    NEXT_QUANTUM
<del>/***********************</del>
void next_quantum (PROCESSOR *nodeP)
 int idx;
 nodeP->quanta idx++;
 if (nodeP->quanta_idx >= NQUANTA)
  nodeP->quanta idx = 0;
                 = nodeP->quanta_idx;
 nodeP->quanta[idx].used
                        = 0;
 nodeP->OSoverhead
                       -= nodeP->quanta[idx].OSportion;
 nodeP->quanta[idx].OSportion = 0;
 nodeP->quanta[idx].actual_load = nodeP->n_local_procs;
 nodeP->quanta[idx].virtual_load = nodeP->n_virtual_procs;
For Dynamic Load
                      Balancing
    CALC_AVE_LOAD
                                                              */
/*********************
float dyn_calc_ave_load(PROCESSOR *nodeP)
 int total_load,
   idx;
   total_load = 0;
   for (idx=0; idx <NQUANTA; idx++)
     total_load += nodeP->quanta[idx].actual load
            + nodeP->quanta[idx].virtual_load;
   return ((float) total load/NQUANTA);
                                                              */
    CALC_AVE_LOAD
                                                              */
/**************
float calc_ave_load(PROCESSOR *nodeP)
 int total_load,
   idx;
   total load = 0;
   for (idx=0; idx < NQUANTA; idx++)
     total_load += nodeP->quanta[idx].actual_load;
   return ((float) total_load/NQUANTA);
```

```
unsigned short exp_xsubi[] = {55213, 10232, 2721};
unsigned short rnd_xsubi[] = {3, 4, 5};
         process_groups[] = {2, 3, 4};
unsigned short rnd_job[]
                          = \{7, 8, 9\};
void job_creation (short int buffer_no)
 extern float load_value;
 extern int njobs,
         total_pe;
  char
          **io_buffer_blk;
  int
  char
           job_name[16];
  double
           rnd_num,
            rnd_grp,
            new rand;
  FILE
           **jobfile blk;
 static double
                t = 0.0;
 static int
                start = 0,
                finish = 0;
 int
                pgsize = PAGE_SIZE;
 extern void setbuf(FILE *, char *);
 extern double log(double);
 extern double erand48(unsigned short []),
              drand48(void);
 if (buffer_no == 0)
       finish = 0;
       t = 0.0;
       start = 0;
    if (njobs < PAGE SIZE)
     pgsize = njobs;
       if (seedfile == NULL)
        Code to randomize jobstreams for each simulation run
        srand48 ((unsigned) time (NULL));
        for (i = 0; i < 3; i++)
              \exp_x subi[i] = (int)(drand48() * 65535);
      }
      else
       {
        fscanf(seedfile, "%hu\n", &exp_xsubi[0]);
        fscanf (seedfile, "%hu\n", &exp_xsubi[1]);
        fscanf(seedfile, "%hu\n", &exp_xsubi[2]);
        };
      fprintf (monitor file,
              "\n *** erand48 seeds are: %d, %d, and %d ***\n",
             exp_xsubi[0], exp_xsubi[1], exp_xsubi[2]);
};
```

```
finish += pgsize;
if (finish > njobs)
       return;
io_buffer_blk = (char **) calloc (total_pe, sizeof(char *));
jobfile_blk = (FILE **) calloc (total_pe, sizeof(FILE *));
/* open job files */
for (i = 0; i < total_pe; i++)
 *(io_buffer_blk + i) = calloc (BUFSIZ, sizeof(char));
 sprintf (job_name, "jobs%d_pg%d", i, buffer_no);
 *(jobfile_blk + i) = fopen(job_name, "wb");
  setbuf(*(jobfile_blk + i), *(io_buffer_blk + i));
};
/* create job streams */
for (; start < finish; start++)
    /* Generate exp. random number. */
    rnd_num = erand48(exp_xsubi);
    new_rand = -log(rnd_num)/((double)(load_value * total_pe));
    /* establish time of creation */
    t = t + new_rand * 1000000 * AVE_PROC_GROUP * AVE_EXEC_TIME;
    /* give job to random mc */
    rnd_grp = erand48(rnd_job);
    rnd_grp = rnd_grp * 3;
    rnd_num = erand48(rnd xsubi);
   rnd_num = rnd_num * total pe;
   fwrite ((const char *) &process_groups[(int)rnd_grp],
         sizeof(process_groups[(int)rnd_grp]),
         1,
         *(jobfile_blk + (int)rnd_num));
   fwrite ((const char *)&t.
         sizeof(t).
         *(jobfile_blk + (int)rnd num));
 }
 /* close job files */
 for (i = 0; i < total_pe; i++)
   cfree (*(io_buffer_blk + i));
  fclose (*(jobfile_blk + i));
 cfree ((char *) jobfile_blk):
 cfree ((char *) io_buffer_blk);
```

```
This version of initb_module on SparcStation One has
                                                                   */
 /*
     been modified with:
                                                                   */
 /*
       a) Routine int_divisible - double/int division &
                                                                   */
         casting of objects.
 /*
       b) Efficient version of all_queues_empty - if any
 /*
         job stream dries up, simulation terminates!
 /*
       c) Data_has_converged routine made more
 /*
         computationally efficient by computing average
       time once only!
 /*
      d) Further efficiency improvement by finding the
 /*
         minimum and maximum run times only.
                                                                   */
 /*
                                                                   */
 /*
     Date of Modification: 5th December 1991
                                                                   */
       #include <stdio.h>
 #include <stdlib.h>
 #include <unistd.h>
 #include <signal.h>
 #include <math.h>
 #include <time.h>
#include <strings.h>
#include "model params.h"
#include "model_types.h"
#include "globvar.h"
CONSTANT, GLOBAL, & LITERAL
                                                                    */
       DEFINITIONS.
/***********************
extern int
                SkipRun;
extern FILE
               *monitor file;
extern NODE_ENTRY
                     *processor table;
extern int
              njobs,
              process_count,
              total_pe;
extern double
                global clock;
extern char
               **io_buffer_blk;
extern void job_creation (short int);
extern double erand48(unsigned short []);
 void
            MsgQ_Dump(MESS_QUEUE),
          Dump_Msg(MSG_FRAME *);
```

```
WRITE_TO_MSGQ
void Write_To_MsgQ(MESS_QUEUE *queue, MSG_FRAME *msgP)
 MSG_FRAME
                  *curr_message,
          *prev_message;
 if (queue->head == NULL)
  queue->head = msgP;
  queue->tail = msgP;
 }
 else
   prev_message = queue->head;
   curr_message = prev_message;
   while (curr_message != NULL)
    if (msgP->time_stamp >= curr_message->time_stamp)
     prev_message = curr_message;
     curr_message = curr_message->next_pack_ptr;
    else break;
   if (curr_message == queue->head)
     msgP->next_pack_ptr = curr_message;
     queue->head = msgP;
   else
    {
     if (curr_message == NULL)
       queue->tail = msgP;
     prev_message->next_pack_ptr = msgP;
     msgP->next_pack_ptr= curr_message;
/* NodeQ_Dump(nodeP); */
```

```
/*
                                                                  */
    READ_FROM_MSGQ
                                                                  */
                                                                  */
MSG_FRAME *Read_From_MsgQ (MESS_QUEUE *queue)
  MSG_FRAME *msg_P;
 msg P = NULL;
 if (queue->head != NULL)
   msg_P
                = queue->head;
   queue->head
                  = queue->head->next_pack_ptr;
   msg_P->next_pack_ptr = NULL;
   if (queue->head == NULL)
     queue->tail = NULL;
 return (msg_P);
void Delete_Cset (PROCESSOR *nP, int id)
int
       i = 0;
while (i < CSET_SIZE)
   if (nP->cSet[i].id < 0)
          break;
   if (nP->cSet[i].id == id)
     while (i < CSET_SIZE - 1)
      nP->cSet[i].id = nP->cSet[i+1].id;
      nP->cSet[i].ld = nP->cSet[i+1].ld;
      i++;
     };
     nP->cSet[i].id = -1;
     nP - cSet[i].Id = 0;
     return;
   else i++;
};
return;
```

```
void Update_Cset ( PROCESSOR *nP, MSG_FRAME *msg, int sort)
   int
            ld,
           wtLd,
           sgn;
   BOOLEAN
                  sorted;
   sorted = FALSE;
   Delete_Cset (nP, msg->sender_id);
   ld = msg->ack_reply;
   if (1d < 0)
     sgn = -1;
  else
     sgn = 1;
   wtLd = ld; /* weighted: (int) ((ld*2)/msg->distance + 0.5); */
   wtLd *= sgn;
     int
             dv[2],
             p, q, i, j;
     if (sort == DESCEND)
       p = 0; q = 1;
     }else
       p = 1; q = 0;
     i = 0;
     while ((i < CSET_SIZE) && (!sorted))
       if (nP->cSet[i].id < 0)
          sorted = TRUE;
       }
       else
        dv[1] = nP -> cSet[i].ld;
        dv[0] = wtLd;
        if (dv[q] < dv[p])
         {
         j = CSET_SIZE - 1;
         while (j > i)
           nP->cSet[j].id = nP->cSet[j-1].id;
           nP->cSet[j].ld = nP->cSet[j-1].ld;
          j--;
          }
         sorted = TRUE;
        } else i++;
       }
     };
     if (sorted)
      nP->cSet[i].id = msg->sender_id;
      nP->cSet[i].ld = wtLd;
     }
#ifdef DEBUG_MSG
 if ((process_count >= MIN_PID)
    && (process_count <= MAX_PID))
```

```
fprintf (monitor_file, "\n\tComms set:\tNode\tLoad\n");
  while (ld < CSET_SIZE)
   fprintf (monitor_file, "\tN%d\t\t%d\t%d\n",
      nP->node_id, nP->cSet[ld].id, nP->cSet[ld].ld);
   ld++;
#endif
  return;
/**********************
                                                           */
                                                           */
/*
    DUMP CSET
void Dump_Cset ( PROCESSOR *pe)
  int ld = 0;
  fprintf (monitor file, "[");
  while ((Id < CSET SIZE) && (pe->cSet[Id].Id >=0))
   fprintf (monitor file, "%d->%d",
       pe->cSet[ld].id, pe->cSet[ld].ld);
  }
  fprintf (monitor file, "]\n\n");
}
/*
                                                           */
    CSET_LOAD_AVG
int Cset_Load_Avg (PROCESSOR *pe)
{ int ld, total = 0, count = 0;
     = (int) pe->threshold;
  while ((count < CSET_SIZE) && (pe->cSet[count].id >=0))
   count++;
   total = total + pe->cSet[count].ld;
  if (count > 0)
   total = total/count;
   if (total > ld)
     return ld;
  if (1d > 2)
    return ld - 1;
  else return ld;
}
```

```
DUMP_MSG
                                                            */
                                                            */
/**********************
 void Dump_Msg (MSG_FRAME *msg_P)
   fprintf (monitor_file,
       "\tSender%d
                  Dest%d
                          distance%d Arr_Time:%f",
       msg_P->sender_id, msg_P->dest_id, msg_P->distance,
       msg_P->time_stamp / ONE_SECOND);
  fprintf \ (monitor\_file, \ ''\ 'n\ 'Type\%d \ ServNo\%d \ Reply\%d\ 'n\ '',
      msg_P->type, msg_P->service_no, msg_P->ack_reply);
}
    MSGQ_DUMP
                                                            */
void MsgQ_Dump (MESS_QUEUE queue)
  MSG_FRAME *msg_P;
 msg_P = queue.head;
 while (msg P != NULL)
  Dump_Msg (msg_P);
  msg_P = msg_P->next_pack_ptr;
 };
}
/*
   NODEQ_DUMP
                                                           */
/****************
void NodeQ_Dump (PROCESSOR *nodeP)
 fprintf (monitor_file, "\nMsgQ State for node%d at time:%f\n",
     nodeP->node_id, nodeP->sys_real_time / ONE_SECOND);
 MsgQ_Dump (nodeP->queue);
```

```
*/
     Dump_R Time
                                                                      */
                                                                      */
void
          Dump_RTime (void)
 int
         Check_for_Zero (int, int),
 float
               rtime;
 PROCESSOR *nodeP;
 for (i = 0; i < total_pe; i++)
   nodeP
           = processor_table[i].node_adr;
   rtime
           = nodeP->cumul_exist_time
            / nodeP->n_deaths;
   fprintf (monitor_file,
       "N:%4d\tMg:%6.2f\tL:%4d\tRt:%6.2f\tTx:%6d Imm:%4d Dth:%6d",
      (float)nodeP->n_migrates,
      nodeP->n_local_procs,
    rtime/ONE_SECOND,
     Check_for_Zero(nodeP->n_tx,nodeP->n_immigs),
     nodeP->n_immigs,
     nodeP->n_deaths);
  fprintf (monitor_file, "\nCommSet:");
  Dump Cset(nodeP);
 fprintf(monitor_file, "\n\n");
}
int Check_for_Zero (int x, int y)
 if (y == 0)
   return x;
 else
   return x/y;
}
```

```
/*
    ERASE_FROM_MSGQ
MSG\_FRAME * Erase\_From\_MsgQ(MESS\_QUEUE * queue,
              int msg type,
              USER_PROCESS *prc)
 MSG_FRAME
                   *last_msg,
           *this_msg;
 last_msg = this_msg = queue->head;
 do
  if (this_msg == NULL)
   return (NULL);
  else
   if (this_msg->service_no == msg_type)
     if ((prc == NULL) || (this_msg->mess_data == prc))
      if (last_msg == this_msg)
       queue->head = this_msg->next_pack_ptr;
       if (this msg == queue->tail)
         queue->tail = queue->head;
       }
      else
       last_msg->next_pack_ptr = this_msg->next_pack_ptr;
     this msg->next_pack_ptr = NULL;
      if (this_msg == queue->tail)
       queue->tail = last msg;
       free ((char *) this_msg);
      return (this_msg);
   last_msg = this_msg;
   this_msg = this_msg->next_pack_ptr;
  } while (TRUE);
```

```
/*
/*
    RESTART_ CLOCK
/*
void restart_clock(PROCESSOR *nodeP,
          int itype,
          double max time,
          USER_PROCESS *prc)
 double
           time stamp;
          *Tim_Text(int);
 char
 BOOLEAN
               reset;
 time_stamp = nodeP->sys_real_time + max_time;
 reset = (!nodeP->timq.itype)
     || ((nodeP->timq.itype != WAITP_TOUT) && (prc == NULL))
     || ((nodeP->timq.itype) && (time_stamp < nodeP->timq.mintime))
     || ((nodeP->timq.itype == WAITP_TOUT) && (itype == WAITP_TOUT)
        && (time_stamp > nodeP->timq.mintime));
 if (reset)
 {
   nodeP->timq.mintime = time_stamp;
   nodeP->timq.itype = itype;
   if (itype == SRQ_TOUT)
     nodeP->timq.minP = prc;
 if (prc != NULL)
   prc->timeout = time_stamp;
   prc->itype = itype;
#ifdef DEBUG MSG
 if ((process_count >= MIN_PID)
    && (process_count <= MAX_PID))
   fprintf (monitor_file,
       "\nNode%d GMT%8.3f \tTripT%4.3f \tTIMER:%s\t Etime:%8.3f\n",
       nodeP->node id,
       nodeP->sys_real_time / ONE_SECOND,
      max time / ONE_SECOND,
      Tim Text(itype),
      time_stamp / ONE_SECOND);
#endif
```

*/

*/

```
*/
/*
   SELECT_NEXT_TIMEOUT
                                                                  */
/*
void Select_next_timeout(PROCESSOR *nodeP)
 USER_PROCESS *proc:
 proc
        = nodeP->firstjb_ptr;
 nodeP->timq.itype = 0;
 nodeP->timq.minP = NULL;
 while (proc != NULL)
  if (proc->timeout)
   nodeP->timq.mintime = proc->timeout;
   nodeP->timq.itype = proc->itype;
   nodeP->timq.minP = proc;
   break;
  else
   proc = proc->nxt_proc_ptr;
 while (proc != NULL)
  if ((proc->timeout) &&
    (proc->timeout < nodeP->timq.mintime))
   nodeP->timq.mintime = proc->timeout;
   nodeP->timq.itype = proc->itype;
   nodeP->timq.minP
                      = proc;
  };
  proc = proc->nxt_proc_ptr;
}
```

```
Tim Text
                                                                        */
/* ==
      *Tim Text(int tcode)
  switch (tcode)
   case RCV_AVG_TOUT:
            return ("RCV_AVG_TOUT\0");
   case HIGH_TOUT: return ("HIGH_TOUT\0");
   case LOW_TOUT : return ("LOW_TOUT\0");
   case SRQ\_TOUT: return ("SRQ\_TOUT\0");
   case PRB_TOUT : return ("PRB_TOUT\0");
   case WAITP_TOUT: return ("WAITP_TOUT\0");
   default: return ("Timer UNDEFINED\0");
  }/* switch */
}
/*
     PAGE FAULTING
/*
     ****************
BOOLEAN
            page_faulting (PROCESSOR *nodeP)
short int
                  node_id;
FILE
                  *tempf_ptr;
char
                  prev fname[16];
/* if (!positive_response) */
   return FAILED;
 node_id = nodeP->node id;
 sprintf (prev_fname, nodeP->event_fname);
 tempf_ptr = nodeP->job_file;
 sprintf (nodeP->event fname, "jobs%d pg%d",
      node_id, nodeP->buffer_no + 1);
 if (init_eventfile (nodeP, node id) == FAILED)
 {
  job_creation (nodeP->buffer no + 1);
  if (init_eventfile (nodeP, node_id) == SUCCESS)
   (nodeP->buffer_no)++;
    fclose (tempf ptr);
    unlink (prev fname);
  }else
     nodeP->job_file = tempf_ptr;
   sprintf (nodeP->event_fname, prev_fname);
 } else
   (nodeP->buffer_no)++;
    fclose (tempf_ptr);
   unlink (prev_fname);
return SUCCESS;
```

```
/*
/*
     INIT_EVENTFILE
                                                  */
/***************
    init_eventfile(PROCESSOR *nodeP, short int node_id)
 int
                  nxt_job;
 double
           nxt_arrival_time;
  if ((nodeP->job_file = fopen(nodeP->event_fname, "rb"))
   == NULL) return FAILED;
  else
   setbuf(nodeP->job_file, *(io_buffer_blk + node_id));
   fread ((char *)&nxt_job, sizeof(nxt_job),
       1,
       nodeP->job file);
   fread ((char *)&nxt_arrival_time, sizeof(nxt_arrival_time),
       1,
       nodeP->job file);
   process_count++;
   nodeP->nxt_arr_time = nxt_arrival_time;
   return SUCCESS;
}
```

```
CLOSE_EVENTFILES
                                                    */
      close_eventfiles (void)
{
 int
              idx;
 PROCESSOR
                     *shadow node;
 for (idx = 0; idx < total_pe; idx++)
  shadow_node = processor_table[idx].node adr;
  fclose(shadow_node->job_file);
  { /* flush process queue */
   USER_PROCESS *process_ptr,
                   *del_prc;
   process_ptr = shadow node->firstjb ptr;
   while (process_ptr != NULL)
    del_prc = process_ptr;
    process_ptr = process_ptr->nxt_proc_ptr;
    free ((char *) del_prc);
   };
  { /* flush message queue */
   MSG_FRAME
                     *msg_P,
                   *del_msg;
    msg_P = shadow_node->queue.head;
    while (msg_P != NULL)
     del_msg = msg_P;
     msg\_P
               = msg P->next pack ptr;
     free ((char *) del_msg);
    };
  }
#ifdef DEBUG_MODEL
  fclose(shadow_node->trace_file);
#endif
  cfree (*(io_buffer_blk + idx));
cfree ((char *) (io_buffer_blk));
```

```
*/
    PROCQ_LIST
                                                           */
                                                           */
void procQ_List (PROCESSOR *nodeP)
 USER PROCESS
                   *proc_P;
 int
 length = 0;
 proc_P = nodeP->firstjb ptr;
 fprintf (monitor_file,"\nNode%d PQ: ", nodeP->node_id);
 while (proc P!= NULL)
   if (length++ % 10);
   else fprintf(monitor_file,"\n");
   fprintf (monitor_file, "pid:%d, ", proc_P->upid);
   proc_P = proc_P->nxt_proc_ptr;
 fprintf(monitor_file,"\n");
   PROCQ_LENGTH
                                                          */
ProcQ_Length (USER_PROCESS *queue)
 int
       length;
 length = 0;
 while (queue != NULL)
   length++;
   queue = queue->nxt_proc_ptr;
 return (length);
/*
                                                          */
/*
   MSG Q_LENGTH
                                                          */
                                                          */
/****************
int MsgQ_Length (MSG_FRAME *queue)
 int
      length;
 length = 0;
 while (queue != NULL)
  length++;
  queue = queue->next_pack_ptr;
 return (length);
```

```
a\,n\,y\,\_\,q\,u\,e\,u\,e\,\,\_\,e\,m\,p\,t\,y
                                               */
                                               */
BOOLEAN any_queue_empty(void)
PROCESSOR
             *nodeP;
         idx:
idx
      = 0;
if (process_count/njobs > 0.60)
 while (idx < total pe)
  nodeP = processor_table[idx].node adr;
  if (feof(nodeP->job_file)
   && (nodeP->firstjb_ptr == NULL))
   return (TRUE);
  idx++;
return (FALSE);
int_divisible
                                               */
BOOLEAN int divisible (double dividend, float divisor)
if ((((int) (dividend / divisor)) * divisor)
 == dividend)
  return TRUE;
else
  return FALSE;
GET TIME
char *Get Time(void)
{
time t clock;
time (&clock);
return (asctime(gmtime(&clock)));
```

```
randmc_id
int randmc_id(PROCESSOR *nodeP)
 int
         mc_id;
 do
 {
 mc_id = (int) (erand48(nodeP->xsubi) * total_pe);
 } while (mc_id == nodeP->node_id);
return (mc_id);
EXIT_DUMP
void exit_dump (PROCESSOR *nodeP, USER_PROCESS *procP)
#ifdef DEBUG_DUMP
 fprintf (nodeP->trace_file, "\n\nProcess died at %.2f sec\n",
    nodeP->sys_real_time / ONE_SECOND);
 fprintf (nodeP->trace_file, "----\n\n");
 fprintf (nodeP->trace_file, "Exec %.2f sec Exist %.2f sec\n",
    procP->exec_time / ONE_SECOND,
    procP->exist_time / ONE_SECOND);
 fprintf (nodeP->trace_file, "Response Ratio = \%.2f sec\n\n\n",
    procP->exec_time / procP->exist_time);
#endif
```

```
/*
     PERFORMANCE_DUMP
                                                                          */
/*
     *************************
void performance_dump( PROCESSOR *nodeP)
 if (nodeP != NULL)
  fprintf (nodeP->trace_file,
       "\nPerformance dump at %.2f sec\n",
       nodeP->sys_real_time / ONE_SECOND);
  fprintf (nodeP->trace_file, "no. of procs %d\n\n",
       nodeP->load);
  if (nodeP->n deaths != 0)
    fprintf (nodeP->trace_file, "Ave RT = \%.2f sec\n",
         (nodeP->cumul_exist_time /ONE_SECOND) /
        nodeP->n_deaths);
  else
    fprintf (nodeP->trace_file, "Ave RT = 0.0 \sec n");
 } /* end of outermost IF */
} /* end of performance dump */
void Int_Trace(int code, ...)
{ int
         i;
 PROCESSOR *nodeP;
 fprintf (monitor file,
      "\n --- SIMULATION DUMP PPC:%d --- @%s \n",
       process_count, Get_Time());
 final_compute();
 Dump RTime();
 for (i = 0; i < total_pe; i++)
  nodeP
           = processor_table[i].node_adr;
  NodeQ_Dump (nodeP);
  procQ List (nodeP);
 fprintf(monitor_file, "\n\n");
 fflush(monitor_file);
 signal (SIGINT, Int Trace);
```

```
void Int_Skip(int code, ...)
{    int    Sec;

fprintf (monitor_file,
        "\n --- SKIPPED SIMULATION PPC:%d --- @%s \n",
        process_count, Get_Time());
    Sec = (int) (global_clock / ONE_SECOND);
    if (Sec > 0)
    {
        final_compute();
        SkipRun = TRUE;
    }
    fflush(monitor_file);
    signal (SIGUSR1, Int_Skip);
}
```

This is the main user-interface module for general menu-interaction, initialisation and initiation of simulation runs.

However, this version differs from the previous (common.src) by abandoning the use of pregenerated, multiple jobfiles and their simulation. "Page Number"/buffer_no is used to simulate data generation through demand_paging.

```
#include <stdio.h>
#include <signal.h>
#include <stdlib.h>
#include <unistd.h>
#include <strings.h>
#include "model params.h"
#include "model types.h"
#include "globvar.h"
extern NODE ENTRY
                   *processor table;
A collection of functions to process a dynamic
    network data structure.
FILE
           *seedfile,
           *trace.
           *monitor_file,
           *configfile,
           *paramfile;
int
            total pe,
           mesh_len,
           positive_response,
           njobs,
           probe_limit,
           lba_index,
           skp_pid;
            load_value,
float
           threshold;
      *menu 1[] = {"1}.
                      EXTEND NETWORK\n",
char
           "2. REDUCE NETWORK\n",
           "3. DISPLAY NETWORK\n",
           "4.
                SIMULATE NETWORK\n",
               EXIT PROGRAM\n\n\n"};
```

DISCONNECT ROW\n",

char

 $*menu_2[] = {"1}.$

```
"2. DELETE ROW\n",
            "3. DELETE NODE\n",
            "4. DELETE ARC\n",
            "9. RETURN TO MAIN\n\n"};
       *menu_3[] = {"1. GENERATE NEW ROW\n",
char
            "2. MESH ROW A & B\n",
            "3. ADD PATH\n",
            "4.
               ADD NODE\n".
            "9. RETURN TO MAIN\n\n"};
main - Test harness for network processing functions.
main (void)
/* declare all local variables.
 int
          menu id,
          new_menu = 1,
          option, batch_mode;
 BOOLEAN
               exit_loop;
 void
           run simulation(void);
 int
          menu_selection(int);
 /* main: process code.
                                   */
 batch_mode = BATCH_RUN;
 if (!batch_mode)
  monitor_file = stdout;
  {/* select option and perform function
   menu_id = new_menu;
   option = menu_selection(menu_id);
   switch (menu_id)
   {/* perform selected function
                                     */
    case 1: /*
        switch (option)
        {/* perform selected function
          case 1: /* */
             new_menu = 3;
             break;
          case 2: /* */
             new menu = 2;
             break;
          case 3: break;
          case 4: run_simulation();
             break;
         default: /* exit from program
                                          */
             break;
```

```
}/* switch */
           break;
      case 2: /*
            switch (option)
            {/* perform selected function
              case 1: /* push a given data item onto stack
                                                               */
              case 2: /* pop a data item off the stack
                   break;
              case 3: /* display the contents of the stack
                                                               */
                   break:
              case 4: break;
             default: /* exit from program
                                                          */
                   new menu = 1;
                   break;
           }/* switch */
           break;
      case 3: /* display the contents of the stack
           switch (option)
            {/* perform selected function
                                                          */
              case 1: /* push a given data item onto stack
                                                               */
              case 2: /* pop a data item off the stack
                   break;
              case 3: /* display the contents of the stack
                   break;
              case 4: break;
             default: /* exit from program
                                                          */
                  new menu = 1;
                  break;
           }/* switch */
          break;
      default:/* exit from program
           break;
    };/* switch */
    exit_loop = (menu_id == 1) && (option == 9);
   } while (!exit loop);
 else
   { char trcName[15];
     skp_pid = getpid();
     sprintf (trcName, "Pid%d.trc", skp_pid);
     monitor_file = fopen (trcName, "w");
     run_simulation();
 fclose (monitor_file);
} /* main */
```

```
menu_selection - function to display menu and return
          index of selected option
              menu_selection(int menu id)
{ /* implementation */
 char
       choice;
 void
       Display_menu(char *[], int);
 switch (menu_id)
  {/* perform selected function
                               */
  case 1: Display_menu(menu_1, 5);
     break;
  case 2: Display_menu(menu_2, 5);
     break;
  case 3: Display_menu(menu_3, 5);
     break;
  default: /* */
     break;
  }/* switch */
 choice = getc(stdin) - 48;
 if ( (choice < 1) || (choice >9) )
  return EOF:
 else return choice;
}/* menu_selection */
Display menu
                                                    */
                                                    */
void Display_menu(char *menu[], int entries)
       index;
 int
 for (index = 0; index < entries; index++)
   printf ("%s\n",menu[index]);
```

```
run_simulation
                                                                           */
void run_simulation (void)
 extern void
                create_config file(int),
                 create_params_file(int).
                 simulate_network(void);
 extern void
                Build Network(),
                write_intro_script();
 BOOLEAN
                positive_response = FALSE;
 void
            Initialise_Netmodel(NODE_ENTRY[]),
            clear_workspace(NODE ENTRY []).
            init_model_params(void);
 char
             *Get Time(void);
 int
            run_count = 1,
            prev total = 0;
 fprintf (monitor file,
      "\n *** SIMULATION COMMENCED ON %s \n\n", Get_Time());
 trace = fopen("main.trc", "a");
 create_config_file (BATCH_RUN);
 create_params_file (BATCH RUN);
 configfile = fopen ("config.dat", "r");
 write_intro_script();
 while (!feof(configfile))
  processor_table = (NODE ENTRY *)
           calloc (total_pe, sizeof(NODE_ENTRY));
  fscanf (configfile, "%d", & mesh_len);
  if (prev_total == total_pe)
   fscanf (configfile, "%d", &positive_response);
  }
  else
    prev_total = total pe;
  Build_Network();
  seedfile = fopen("seedfile.dat", "r");
  paramfile = fopen ("pramfile.dat", "r");
  fscanf (paramfile, "%f", &load value);
  while (!feof(paramfile))
   fprintf (monitor_file,
        "\nSIMULATION no.%d STARTED ON %s",
        run count, Get Time());
   Initialise_Netmodel(processor_table);
   init model params();
   signal (SIGINT, Int_Trace);
   signal (SIGUSR1, Int_Skip);
   simulate network();
   fprintf (monitor file,
        "____\n\n\014");
```

```
fflush (monitor_file);
   fflush (trace);
   run_count++;
   fscanf (paramfile, "%f", &load_value);
  fclose (paramfile);
  fclose (seedfile);
  clear_workspace(processor_table);
  write_intro_script();
 fclose(configfile);
 fclose(trace);
}
    clear_workspace
                                                                    */
/*
                                                                    */
void clear_workspace (NODE_ENTRY processor_table[])
  int
         idx;
  for (idx = 0; idx < total_pe; idx++)
    cfree ((char *) processor table[idx].node adr->links);
    free ((char *) processor_table[idx].node_adr);
  cfree ((char *) processor_table);
```

```
Initialise Netmodel
                                                                           */
void Initialise_Netmodel(NODE_ENTRY processor_table[])
{ /* implementation */
 PROCESSOR
                *nodeP;
 extern double drand48(void):
          idx1, idx2;
 for (idx1 = 0; idx1 < total pe; idx1++)
    nodeP = processor_table[idx1].node_adr;
   /* Construct data node */
   processor_table[idx1].await_sync = TRUE;
   nodeP->load
                       = N_SYS_PROCS;
                        = 0:
   nodeP->pwait set
   nodeP->low t set
                         = 0:
   nodeP->high t set
   nodeP->rplimit
                       = 0;
   nodeP->n_deaths
                        = 0;
   nodeP->n_local_procs = N_SYS_PROCS;
   nodeP->n_virtual_procs = 0;
   nodeP->n tx
   nodeP->n_rx
                       = 0;
                         = 0;
   nodeP->n migrates
   nodeP->n_immigs
                         = 0;
   nodeP->cumul_exist_time = 0.0;
   nodeP->sys real time = 0.0;
   nodeP->last_perf_dump_time = 0.0;
   nodeP->next elapsed second = 1000000.0;
   nodeP->n_active_local_procs = N SYS PROCS:
   nodeP->quanta idx
                         = 0;
   nodeP->OSoverhead
                          = 0;
   for (idx2 = 0; idx2 < NQUANTA; idx2++)
     nodeP->quanta[idx2].actual_load = N_SYS_PROCS;
     nodeP->quanta[idx2].virtual load = 0;
     nodeP->quanta[idx2].OSportion = 0;
     nodeP->quanta[idx2].used
    for (idx2 = 0; idx2 < 3; idx2++)
    {
       nodeP->xsubi[idx2] = (int) (drand48() * total_pe);
     }while (nodeP->xsubi[idx2] == nodeP->node id);
    for (idx2 = 0; idx2 < CSET_SIZE; idx2++)
      nodeP -> cSet[idx2].id = -1;
      nodeP -> cSet[idx2].Id = 0;
   nodeP->buffer_no = 0;
   nodeP->curr_process= NULL:
   nodeP->firstjb_ptr = NULL:
   nodeP->lastjb_ptr = NULL:
   nodeP->queue.tail = NULL:
   nodeP->queue.head = NULL:
   nodeP->timq.minP = NULL;
```

```
nodeP->timq.itype = 0;
   nodeP->timq.mintime = 0;
}
/* _____*/
     init_model_params
/*
                                                                       */
void init model params(void)
{
 extern void job_creation(short int);
         *lba_text(void);
 char
 int
         idx;
   fscanf(paramfile, "%d", &njobs);
   fscanf (paramfile, "%d", &positive response);
   fscanf(paramfile, "%d", &lba_index);
   if (positive_response)
   {
        if (system ("rm jobs*"));
        job creation (0);
   }
   fprintf (monitor file, "\n%15s%9d\t", "No. of Processors = ", total_pe);
                                                 = ", load_value);
   fprintf (monitor_file, "%15s%1.2f", "Load value
   fprintf (monitor_file, "\n%15s%9d\t", "Total No. of jobs = ", njobs);
   fprintf (monitor_file, "%15s%s\n", "Ld Balancing Alg. = ", lba_text());
   fscanf (paramfile, "%f", &threshold);
   for (idx = 0; idx < total_pe; idx++)
    (processor_table[idx].node_adr)->threshold = (int) threshold;
   if (lba_index > 0)
    fscanf (paramfile, "%d", &probe_limit);
    fprintf (monitor file, "\n%15s%9.2f\t",
              "Threshold Level = ", threshold);
    switch (lba index)
    {/* perform selected function */
     case 2: fprintf (monitor_file, "%15s%9d",
                       = ", probe_limit);
          "Probe Limit
          break:
     case 3:
     case 4:
     case 5:
     case 6:
     case 7:
     case 8:
      case 9:
      case 10: fprintf (monitor_file, "%20s%2d",
                "Distance
                             = ", probe_limit);
      default: break;
    }
   fprintf (monitor_file, "\n%15s%1.2f%%\n", "Convergence to < ",
                 100 * CONVERGE_FACTOR);
   fprintf (monitor_file, "\n");
 }
```

```
*lba_text(void)
switch (lba_index)
{/* perform selected function
 case 0: return ("No Load Balancing\0");
 case 1: return ("Random L/B Policy\0");
 case 2: return ("Threshold L/B Policy\0");
case 3: return ("Reverse Broadcast\0");
 case 4: return ("Threshold Broadcast\0");
case 5: return ("Threshold Nbor Policy\0");
case 6: return ("Adaptive Th. Bcast\0");
case 7: return ("Adaptive Rev. Bcast\0");
case 8: return ("Global Avg Nbor Policy\0");
case 9: return ("Global Avg Broadcast\0");
case 10: return ("Global Avg Rcv Bcast\0");
default: return ("L/B policy UNDEFINED\0");
}/* switch */
```