

Ramsey numbers of cubes versus cliques

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Abstract

The cube graph Q_n is the skeleton of the n -dimensional cube. It is an n -regular graph on 2^n vertices. The Ramsey number $r(Q_n, K_s)$ is the minimum N such that every graph of order N contains the cube graph Q_n or an independent set of order s . In 1983, Burr and Erdős asked whether the simple lower bound $r(Q_n, K_s) \geq (s-1)(2^n-1)+1$ is tight for s fixed and n sufficiently large. We make progress on this problem, obtaining the first upper bound which is within a constant factor of the lower bound.

1 Introduction

For graphs G and H , the *Ramsey number* $r(G, H)$ is defined to be the smallest natural number N such that every red/blue edge-coloring of the complete graph K_N on N vertices contains a red copy of G or a blue copy of H .

One obvious construction, noted by Chvátal and Harary [16], which gives a lower bound for these numbers is to take $\chi(H) - 1$ disjoint red cliques of size $|G| - 1$ and to connect every pair of vertices which are in different cliques by a blue edge. If G is connected, the resulting graph contains neither a red copy of G nor a blue copy of H , so that $r(G, H) \geq (|G| - 1)(\chi(H) - 1) + 1$. Burr [8] strengthened this bound by noting that if $\sigma(H)$ is the smallest color class in any $\chi(H)$ -coloring of the vertices of H , we may add a further red clique of size $\sigma(H) - 1$, obtaining

$$r(G, H) \geq (|G| - 1)(\chi(H) - 1) + \sigma(H).$$

Following Burr and Erdős [8, 11], we say that a graph G is *H-good* if the Ramsey number $r(G, H)$ is equal to this bound. If \mathcal{G} is a family of graphs, we say that \mathcal{G} is *H-good* if all sufficiently large graphs in \mathcal{G} are *H-good*. When $H = K_s$, where $\sigma(K_s) = 1$, we simply say that G or \mathcal{G} is *s-good*.

The classical result on Ramsey goodness, which predates the definition, is the theorem of Chvátal [15] showing that all trees are *s-good* for any s . On the other hand, the family of trees is not *H-good* for every graph H . For example [13], a construction of $K_{2,2}$ -free graphs due to Brown [7] allows one to show that there is a constant $c < \frac{1}{2}$ such that

$$r(K_{1,t}, K_{2,2}) \geq t + \sqrt{t} - t^c$$

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for t sufficiently large. This is clearly larger than $(|K_{1,t}| - 1)(\chi(K_{2,2}) - 1) + \sigma(K_{2,2}) = t + 2$.

In an effort to determine what properties contribute to being Ramsey good, Burr and Erdős [9, 11] conjectured that if Δ is fixed then the family of graphs with bounded maximum degree Δ is s -good for any s (and perhaps even H -good for all H). This conjecture holds for bipartite graphs H [12] but is false in general, as shown by Brandt [6]. He proved that for $\Delta \geq \Delta_0$ almost every Δ -regular graph on a sufficiently large number of vertices is not even 3-good. His result (and a similar result in [30]) actually proves something stronger, namely, that if a graph G has strong expansion properties then it cannot be 3-good.

On the other hand, it has been shown that if a family of graphs exhibits poor expansion properties then it will tend to be good [1, 30]. To state the relevant results, we define the *bandwidth* of a graph G to be the smallest number ℓ for which there exists an ordering v_1, \dots, v_n of the vertices of G such that every edge $v_i v_j$ satisfies $|i - j| \leq \ell$. This parameter is known to be intimately linked to the expansion properties of the graph. In particular, any bounded-degree graph with poor expansion properties will have sublinear bandwidth [4].

The first such result, shown by Burr and Erdős [11], states that for any fixed ℓ the family of connected graphs with bandwidth at most ℓ is s -good for any s . This result was recently extended by Allen, Brightwell and Skokan [1], who showed that the set of connected graphs with bandwidth at most ℓ is H -good for every H . Their result even allows the bandwidth ℓ to grow at a reasonable rate with the size of the graph G . If G is known to have bounded maximum degree, their results are particularly strong, saying that for any Δ and any fixed graph H there exists a constant c such that if G is a graph on n vertices with maximum degree Δ and bandwidth at most cn then G is H -good.

Many of the original problems of Burr and Erdős [11] have now been resolved [30] but one that remains open is to determine whether the family of hypercubes is s -good for every s . The *hypercube* Q_n is the graph on vertex set $\{0, 1\}^n$ where two vertices are connected by an edge if and only if they differ in exactly one coordinate. This family of graphs has sublinear bandwidth but does not have bounded degree, so the result of Allen, Brightwell and Skokan does not apply.

To get a first bound for $r(Q_n, K_3)$, note that a simple greedy embedding implies that any graph with maximum degree d and at least $dn + 2^n$ vertices has a copy of Q_n in its complement. Suppose now that the edges of a complete graph have been 2-colored in red and blue and there is neither a blue triangle nor a red copy of Q_n . Then, since the blue neighborhood of any vertex forms a red clique, the maximum degree in blue is at most $2^n - 1$. Hence, the graph must have at most $(2^n - 1)n + 2^n < 2^n(n + 1)$ vertices. We may therefore conclude that $r(Q_n, K_3) \leq 2^n(n + 1)$.

It is not hard to extend this argument to show that for any s there exists a constant c_s such that $r(Q_n, K_s) \leq c_s 2^n n^{s-2}$. This is essentially the best known bound. Here we improve this bound, obtaining the first upper bound which is within a constant factor of the lower bound.

Theorem 1.1. *For any natural number $s \geq 3$, there exists a constant c_s such that*

$$r(Q_n, K_s) \leq c_s 2^n.$$

The original question of Burr and Erdős [11] relates to s -goodness but it is natural to also ask whether the family of cubes is H -good for any H . For bipartite H , this follows directly from a result of Burr, Erdős, Faudree, Rousseau and Schelp [12]. Our result clearly implies that for any H , there is a constant c_H such that $r(Q_n, H) \leq c_H 2^n$.

For triangles, the rough idea of the proof is to show that if a red/blue edge-coloring of K_N does

not contain a blue triangle then it may be tiled with a collection of red cliques which have low blue density between them. A red copy of Q_n may then be found by inserting subcubes into each of the red cliques and patching them together using the fact that the blue density between these different cliques is low. Throughout the argument, it is very important to keep close control on the size and number of the red cliques as the definition of low blue density in a particular bipartite subgraph will depend on the size of the two cliques forming its endpoints.

For K_4 , the method is an extension of this idea, except the tiling will now consist of cliques which are free of blue triangles, again with low blue density between different cliques. In order to make this useful for embedding cubes, we then have to perform a second level of tiling, splitting each such clique into red subcliques with low density between them, as was already done for triangles. In general, when we consider K_s , there will be $s - 2$ levels of tiling to keep track of. This makes the bookkeeping somewhat complex. Accordingly, we have chosen to present the proof in stages, first considering triangles, then K_4 and only then the general case. Although this leads to some redundancies, it allows us to introduce the additional concepts needed for each step at a reasonable pace.

More generally, one may ask which families of graphs are s -good for all s . A powerful result proved by Nikiforov and Rousseau [30] shows that graphs with small separators are s -good. They used this result to resolve a number of the original questions of Burr and Erdős [11] regarding Ramsey goodness. Let the *degeneracy* $d(G)$ of a graph G be the smallest natural number d such that every induced subgraph of G has a vertex of degree at most d . Furthermore, we say that a graph G has a (t, η) -separator if there exists a vertex subset $T \subseteq V(G)$ such that $|T| \leq t$ and every connected component of $V(G) \setminus T$ has size at most $\eta|V(G)|$. The result of Nikiforov and Rousseau then says that for any $s \geq 3$, $d \geq 1$ and $0 < \gamma < 1$, there exists $\eta > 0$ such that the class \mathcal{G} of d -degenerate graphs G with a $(|V(G)|^{1-\gamma}, \eta)$ -separator is s -good.

We will apply this theorem, together with the Alon-Seymour-Thomas separator theorem for graphs with a forbidden minor [2], to show that for any $s \geq 3$ any family of graphs with a forbidden minor is s -good. A graph H is said to be a *minor* of G if H can be obtained from a subgraph of G by contracting edges. By an H -minor of G , we mean a minor of G which is isomorphic to H . For a graph H , let \mathcal{G}_H be the family of graphs which do not contain an H -minor.

Theorem 1.2. *For every fixed graph H , the class \mathcal{G}_H of graphs G which do not contain an H -minor is s -good for all $s \geq 3$.*

In particular, since the family of planar graphs consists exactly of those graphs which do not contain K_5 or $K_{3,3}$ as a minor, we have the following corollary.

Corollary 1.1. *The family of planar graphs is s -good for all $s \geq 3$.*

A *minor-closed* family \mathcal{G} is a collection of graphs which is closed under taking minors. The graph minor theorem of Robertson and Seymour [31] states that any minor-closed family of graphs may be characterized by a finite collection of forbidden minors. We say that a minor-closed family is nontrivial if it is not the class consisting of all graphs. Note that the following corollary is an immediate consequence of Theorem 1.2 as any nontrivial minor-closed family \mathcal{G} is a subfamily of \mathcal{G}_H , where H is a graph not in \mathcal{G} .

Corollary 1.2. *Any nontrivial minor-closed family of graphs \mathcal{G} is s -good for all $s \geq 3$.*

We will begin, in Section 2, by studying the Ramsey number of cubes versus triangles. We will then show, in Section 3, how our arguments extend to K_4 before treating the general case in Section 4. In Section 5, we will prove Theorem 1.2. We will conclude with some further remarks. All logarithms are base 2 unless otherwise indicated. For the sake of clarity of presentation, we systematically omit floor and ceiling signs whenever they are not crucial. We also do not make any serious attempt to optimize absolute constants in our statements and proofs.

2 Triangle versus cube

The argument works for $n \geq 6$. Consider a coloring of the edges of the complete graph K_N on the vertex set $[N] = \{1, 2, \dots, N\}$ for $N \geq 7000 \cdot 2^n$ with two colors, red and blue, and assume that there are no blue triangles. We will prove that this coloring contains a red Q_n .

2.1 Preprocessing the coloring

For each $d = 0, 1, 2, \dots, \log n + 3$, we use the following procedure to construct a family \mathcal{S} of subsets of $[N]$ (note that $\log n + 3 \leq n$ for $n \geq 6$):

If there exists a set S which induces a red clique of order exactly $4 \cdot 2^{n-d}$, then arbitrarily choose one, add it to the family \mathcal{S} , and remove the vertices of the clique from $[N]$. We define the *codimension* $d(S)$ of such a set as $d(S) = d$. When there are no more such red cliques, continue to the next value of d . In the end, if we have $\sum_{S \in \mathcal{S}} |S| \geq \frac{N}{2}$, then we let \mathcal{S} be our family. Otherwise, if $\sum_{S \in \mathcal{S}} |S| < \frac{N}{2}$, we add the set of remaining vertices to \mathcal{S} , and declare it to have codimension zero (note that this set has size at least $\frac{N}{2} \geq 4 \cdot 2^n$).

In either of the cases, we have $\sum_{S \in \mathcal{S}} |S| \geq \frac{N}{2}$. Note that we also have the following properties.

Proposition 2.1. (i) For an integer $i \geq 1$, let $X = \bigcup_{S \in \mathcal{S}, d(S) \geq i} S$. Then each vertex $v \in [N]$ has at most 2^{n-i+3} blue neighbors in X .
(ii) For every set $S \in \mathcal{S}$, the subgraph induced by S has maximum blue degree at most $\frac{2^{n-d(S)}}{2n}$.

Proof. (i) Given $i \geq 1$ and a set X defined as above, note that no subset of X induces a red clique of size at least $4 \cdot 2^{n-i+1} = 2^{n-i+3}$, since such a set would have been added to \mathcal{S} in the previous round. On the other hand, since there are no blue triangles, the blue neighborhood of every vertex induces a red clique. Therefore, for every $v \in [N]$, v has at most 2^{n-i+3} blue neighbors in X .

(ii) The claim is trivially true for $d > 0$, since every set of codimension at least 1 is a red clique. For $d = 0$, if S is not a red clique, then it is the set of codimension zero added to \mathcal{S} after the sets of codimensions 0 to $\log n + 3$. Thus, there are no subsets of S which induce a red clique of size at least $4 \cdot 2^{n-(\log n + 3)}$, since such a set would have been added earlier to the family \mathcal{S} . On the other hand, since there are no blue triangles, the blue neighborhood of every vertex induces a red clique. Therefore, every vertex of S has at most $4 \cdot 2^{n-(\log n + 3)} \leq \frac{2^n}{2n}$ blue neighbors in S . \square

2.2 Tiling the cube

Our strategy is to first decompose the cube Q_n into smaller cubes and then to embed it piece by piece into K_N , placing one of the subcubes in each of the subsets from \mathcal{S} . We represent the subcubes

of Q_n by using vectors in $\{0, 1, *\}^n$. For example, the vector $(0, 1, *, \dots, *)$ will represent the subcube $\{(0, 1, x_1, \dots, x_{n-2}) : (x_1, \dots, x_{n-2}) \in \{0, 1\}^{n-2}\}$.

We seek a *tiling* of Q_n , which we define as a collection of vertex-disjoint cubes that covers all the vertices of Q_n and uses only cubes that have all their fixed coordinates at the start of the vector. That is, they are of the form $C = (a_1, a_2, \dots, a_d, *, *, \dots, *)$ for some $d \geq 0$ and $a_1, a_2, \dots, a_d \in \{0, 1\}$. We will call such a cube *special* and let d be the *codimension* $d(C)$ of C . We say that two disjoint cubes C and C' are *adjacent*, if there are vertices $x \in C$ and $x' \in C'$ which are adjacent in Q_n . Note that the following list of properties holds for such cubes and a tiling composed from them.

Proposition 2.2. (i) *If two special cubes C and C' intersect, then we have $C \subset C'$ or $C' \subset C$.*

(ii) *Two disjoint cubes $C = (a_1, \dots, a_d, *, \dots, *)$ and $C' = (b_1, \dots, b_{d'}, *, \dots, *)$ for $d' \leq d$ are adjacent if and only if $(a_1, \dots, a_{d'})$ and $(b_1, \dots, b_{d'})$ differ in exactly one coordinate.*

(iii) *For a tiling \mathcal{C} and a special cube $C \in \mathcal{C}$ of codimension d , there are at most d other special cubes $C' \in \mathcal{C}$ of codimension $d(C') \leq d$ which are adjacent to C .*

Proof. (i) Let $d = d(C)$, $d' = d(C')$ and, without loss of generality, suppose that $d' \leq d$. Suppose that $(a_1, \dots, a_n) \in C \cap C'$. Then, since we are only considering special cubes, we must have $C = (a_1, \dots, a_d, *, \dots, *)$ and $C' = (a_1, \dots, a_{d'}, *, \dots, *)$. However, we then have $C \subset C'$.

(ii) If $(a_1, \dots, a_{d'}) = (b_1, \dots, b_{d'})$, then we have $C \subset C'$, and thus we may assume that this is not the case. Cubes C and C' are adjacent if and only if there are two vectors $\mathbf{v} = (a_1, \dots, a_d, x_{d+1}, \dots, x_n)$ and $\mathbf{w} = (b_1, \dots, b_{d'}, y_{d'+1}, \dots, y_n)$ in $\{0, 1\}^n$ which differ in exactly one coordinate. However, since the two vectors restricted to the first d' coordinates are already different, this can happen only if $(a_1, \dots, a_{d'})$ and $(b_1, \dots, b_{d'})$ differ in exactly one coordinate. Moreover, one can see that if this is the case, then there is an assignment of values to x_i and y_j so that \mathbf{v} and \mathbf{w} indeed differ in exactly one coordinate.

(iii) Since C has codimension d , there are only d coordinates that one can ‘flip’ from C to obtain a cube adjacent to C . \square

As we have already mentioned, our tiling \mathcal{C} of the cube Q_n will be constructed in correspondence with the family \mathcal{S} constructed in Section 2.1. We will construct the tiling \mathcal{C} by finding cubes of the tiling one by one. We slightly abuse notation and use \mathcal{C} also to denote the ‘partial’ tiling, where only part of the cube Q_n is covered. At each step, we will find a subcube C which covers some non-covered part of Q_n and assign it to some set $S_C \in \mathcal{S}$. We say that such an assignment is *proper* if the following properties hold.

Proper assignment.

1. If C has codimension d , then S_C has codimension d .
2. Suppose that C is adjacent to some cube C' already in the tiling and that C' is assigned to $S_{C'}$. Then the bipartite graph induced by S_C and $S_{C'}$ contains at most $\frac{|S_C||S_{C'}|}{16\delta^2}$ blue edges, where $\delta = \max\{d(C), d(C')\}$.

Our algorithm for finding the tiling \mathcal{C} and the corresponding sets in \mathcal{S} is as follows.

Tiling Algorithm. At each step, consider all possible special cubes C which

- (a) are disjoint from the cubes $C' \in \mathcal{C}$,
- (b) have $d(C) \geq d(C')$ for all $C' \in \mathcal{C}$, and
- (c) for which there exists a set $S \in \mathcal{S}$ which has not yet been assigned and such that assigning C to S is a proper assignment.

Take a cube C_0 of minimum codimension satisfying (a), (b) and (c) and add it to the tiling. Assign C_0 to the set $S_{C_0} \in \mathcal{S}$ given by (c).

The following proposition shows that the algorithm will terminate successfully.

Proposition 2.3. *If the tiling is not complete, then the algorithm always chooses a cube from a non-empty collection.*

Proof. If \mathcal{S} contains a set of codimension zero, then the algorithm will choose the whole cube Q_n in the first step and assign it to a set of codimension zero. In this case we have a trivial tiling and there is nothing to prove.

Thus, we may assume that \mathcal{S} contains no set of codimension zero. If this is the case, note that for all $C \in \mathcal{C}$, we have $|S_C| = 4|C|$. Let \mathcal{S}'' be the subfamily of \mathcal{S} consisting of sets which are already assigned to some cube in \mathcal{C} and let $\mathcal{S}' = \mathcal{S} \setminus \mathcal{S}''$. Note that

$$\left| \bigcup_{S \in \mathcal{S}'} S \right| = \left| \bigcup_{S \in \mathcal{S}} S \right| - \left| \bigcup_{S \in \mathcal{S}''} S \right| \geq \frac{N}{2} - \left| \bigcup_{C \in \mathcal{C}} S_C \right| \geq 3500 \cdot 2^n - 4 \cdot 2^n \geq 3496 \cdot 2^n.$$

For each i , let $\mathcal{S}'_i = \{S \in \mathcal{S}' : d(S) = i\}$. If $|\mathcal{S}'_i| \leq 32i^3$ for all i , then we have

$$\left| \bigcup_{S \in \mathcal{S}'} S \right| = \sum_i \left| \bigcup_{S \in \mathcal{S}'_i} S \right| \leq \sum_i 32i^3 \cdot 4 \cdot 2^{n-i} = 128 \cdot 2^n \sum_i \frac{i^3}{2^i} = 3328 \cdot 2^n,$$

which is a contradiction (we used the fact that $\sum_i \frac{i^3}{2^i} = 26$). Therefore, there exists an index i for which $|\mathcal{S}'_i| > 32i^3$. We have $i > 0$ since \mathcal{S} contains no set of codimension zero.

Since the tiling is not complete, there exists a vertex $(a_1, \dots, a_n) \in Q_n$ which is not yet covered by any of the cubes already in \mathcal{C} . Consider the cube $C = (a_1, \dots, a_i, *, \dots, *)$. By Proposition 2.2, there are at most i cubes in our partial tiling \mathcal{C} of codimension at most i which are adjacent to this cube. Let C_1, \dots, C_j for $j \leq i$ be these cubes and fix an index $a \leq j$. We say that a set $S \in \mathcal{S}'_i$ is *bad* for S_{C_a} if there are at least $\frac{1}{16i^2} |S_{C_a}| |S|$ blue edges between S_{C_a} and S . We claim that there are at most $32i^2$ sets in \mathcal{S}'_i which are bad for S_{C_a} .

Let $X_i = \bigcup_{S \in \mathcal{S}'_i} S$ and note, by Proposition 2.1(i), that there are at most $|S_{C_a}| \cdot 2^{n-i+3}$ blue edges between the sets S_{C_a} and X_i . Each set $S \in \mathcal{S}'_i$ which is bad for S_{C_a} accounts for at least $\frac{1}{16i^2} |S_{C_a}| |S| = \frac{2^{n-i+3}}{32i^2} |S_{C_a}|$ such blue edges (note that $|S| = 4 \cdot 2^{n-i}$). Therefore, in total, there are at most $32i^2$ sets in \mathcal{S}'_i which are in bad relation with S_{C_a} , as claimed above. Since $|\mathcal{S}'_i| > 32i^3$, there exists a set $S \in \mathcal{S}'_i$ which is not bad for any of the sets S_{C_a} for $1 \leq a \leq j \leq i$.

Suppose that in the previous step we embedded some cube of codimension d . In order to show that C satisfies (a) and (b), it suffices to verify that $i \geq d$, since $i \geq d$ implies the fact that C is disjoint from all the other cubes in \mathcal{C} (if C intersects some other cube, then that cube must contain C and therefore also contains (a_1, \dots, a_n) by Proposition 2.2). Furthermore, if this is the case, assigning C

to S is a proper assignment (thus we have (c)), since C_1, \dots, C_j are the only cubes adjacent to C already in the tiling.

Suppose, for the sake of contradiction, that $i < d$ and consider the time t immediately after we last embedded a cube of codimension at most i . At time t , since (a_1, \dots, a_n) was not covered, the cubes in \mathcal{C} are disjoint from C . Moreover, the set of cubes of codimension at most i which are adjacent to C is the same as at the current time. Therefore, C could have been added to the tiling at time t as well, contradicting the fact that we always choose a cube of minimum codimension. Thus we have $i \geq d$, as claimed. \square

Note that as an outcome of our algorithm, we obtain a tiling \mathcal{C} for which every pair of adjacent cubes $C, C' \in \mathcal{C}$ have at most $\frac{1}{16\delta^2}|S_C||S_{C'}|$ blue edges between S_C and $S_{C'}$, where $\delta = \max\{d(C), d(C')\}$.

2.3 Imposing a maximum degree condition

Having constructed the tiling \mathcal{C} and made the assignment of the cubes to sets in \mathcal{S} , we now wish to impose certain maximum degree conditions between the sets S_C for $C \in \mathcal{C}$. For a set $C \in \mathcal{C}$ of codimension $d = d(C)$, let C_1, \dots, C_j be the cubes of codimension at most d which are adjacent to C . Note that we have $j \leq d$ by Proposition 2.2. Moreover, there are at most $\frac{1}{16d^2}|S_C||S_{C_a}|$ blue edges between S_C and S_{C_a} for all $a \leq j$.

Now, for each a , remove all the vertices in S_C which have at least $\frac{1}{8d}|S_{C_a}|$ blue neighbors in S_{C_a} and let T_C be the subset of S_C left after these removals. Since there are at most $\frac{1}{16d^2}|S_C||S_{C_a}|$ blue edges between S_C and S_{C_a} , for each index a , we remove at most $\frac{|S_C|}{2d}$ vertices from S_C , and thus the resulting set T_C is of size at least $\frac{|S_C|}{2} \geq 2 \cdot 2^{n-d}$. Note that all the vertices in T_C have blue degree at most $\frac{1}{8d}|S_{C_a}| \leq \frac{1}{4d}|T_{C_a}|$ in the set T_{C_a} . That is, we have the following property.

Maximum degree condition. For each pair of adjacent cubes $C, C' \in \mathcal{C}$ with $d(C) \geq d(C')$, every vertex in T_C has at most $\frac{1}{4d(C)}|T_{C'}|$ blue neighbors in the set $T_{C'}$.

2.4 Embedding the cube

We now show how to embed Q_n . Recall that Q_n was tiled by cubes in \mathcal{C} , each corresponding to a subset from the family \mathcal{S} . We will greedily embed these cubes one by one into their assigned sets from the family \mathcal{S} , in decreasing order of their codimensions. If there are several cubes of the same codimension, then we arbitrary choose the order between them.

For each $C \in \mathcal{C}$, we will greedily embed the vertices of C into the set $T_C \subseteq S_C$. Let $d = d(C)$. Suppose that we are about to embed $x \in C$ and let $f : Q_n \rightarrow [N]$ denote the partial embedding of the cube Q_n obtained so far. Note that x has at most d neighbors x_1, x_2, \dots, x_j (for $j \leq d$) which are already embedded and belong to a cube other than C . Since we have only embedded cubes of codimension at least d , the maximum degree condition imposed in Section 2.3 implies that all the vertices $f(x_i)$ have blue degree at most $\frac{1}{4d}|T_C|$ in the set T_C . Together, these neighbors forbid at most $\frac{1}{4}|T_C|$ vertices of T_C from being the image of x .

In addition, x has at most $n - d$ neighbors y_1, \dots, y_k (for $k \leq n - d$) which are already embedded and belong to C . By Proposition 2.1, each vertex $f(y_i)$ has blue degree at most $\frac{2^{n-d}}{2^n}$ in the set T_C . Together, these neighbors forbid at most $\frac{2^{n-d}}{2}$ vertices of T_C from being the image of x . Finally,

there are at most $2^{n-d} - 1$ vertices in T_C which are images of some other vertex of C that is already embedded. Therefore, the number of vertices in T_C into which we cannot embed x is at most

$$\frac{1}{4}|T_C| + \frac{1}{2} \cdot 2^{n-d} + (2^{n-d} - 1),$$

which is less than $|T_C|$ since $|T_C| \geq 2 \cdot 2^{n-d}$. Hence, there exists a vertex in T_C which we can choose as an image of x to extend the current partial embedding of the cube. Repeating this procedure until we embed the whole cube Q_n completes the proof.

3 Cliques of order 4

The argument for general cliques is similar to that for triangles given in the previous section. However, there are several new concepts involved. To slowly develop the necessary concepts, we first provide a proof of the next case, which is K_4 versus a cube.

Recall that in the triangle case, we started by iteratively finding sets which formed a red clique (see Section 2.1). Red cliques were a natural choice, since the blue neighborhood of every vertex formed a red clique. Either we were able to find large red cliques or we were able to restrict the maximum blue degree of vertices in some way (see Proposition 2.1). If one attempts to employ the same strategy for the K_4 case, then the natural choice of sets that we should take instead of red cliques are blue triangle-free sets, since the blue neighborhood of every vertex now forms a blue triangle-free set. Suppose that we have found a family \mathcal{S}_1 of blue triangle-free sets. Since blue K_3 -free sets are not as powerful as red cliques for embedding subgraphs, we repeat the whole argument within each blue triangle-free set $S \in \mathcal{S}_1$, to obtain red cliques which are subsets of S . By so doing, we obtain a second family of sets \mathcal{S}_2 , consisting of red cliques. We refer to sets in \mathcal{S}_ℓ as *level ℓ sets* and, for a set $S \in \mathcal{S}_\ell$, we define its level as $\ell(S) = \ell$.

In order to find an embedding of the cube Q_n using a strategy similar to that used in the triangle case, we wish to find a tiling of Q_n and, for each cube C in the tiling, a red clique $S_C \in \mathcal{S}_2$ so that for two adjacent cubes C and C' , the sets S_C and $S_{C'}$ stand in ‘good’ relation. However, directly finding such a tiling and an assignment is somewhat difficult since we do not have good control on the blue edges between the sets in \mathcal{S}_2 . To be more precise, suppose that we are given $S_1, S'_1 \in \mathcal{S}_1$ and subsets $S_2 \subset S_1, S'_2 \subset S'_1$ in \mathcal{S}_2 . Then the control on the blue edges between S_2 and S'_2 is ‘inherited’ from the control on the blue edges between S_1 and S'_1 , and thus depends on the relative sizes of S_1 and S'_1 , not on the relative sizes of S_2 and S'_2 as in the triangle case (unless $S_1 = S'_1$). To circumvent this difficulty, we will need to maintain tight control on the edge density between different sets in \mathcal{S}_1 as well as those in \mathcal{S}_2 .

We seek a *double tiling*, which is defined to be a pair $\mathcal{C} = \mathcal{C}_1 \cup \mathcal{C}_2$ of tilings satisfying the property that for every $C_2 \in \mathcal{C}_2$ there exists a cube $C_1 \in \mathcal{C}_1$ such that $C_2 \subset C_1$ (in other words, \mathcal{C}_2 is a *refined* tiling of \mathcal{C}_1). We refer to cubes in \mathcal{C}_ℓ as *level ℓ cubes* and, for a cube $C \in \mathcal{C}_\ell$, we define its level as $\ell(C) = \ell$. Our goal is to find, for each $\ell = 1, 2$, an assignment of cubes in \mathcal{C}_ℓ to sets in \mathcal{S}_ℓ . The following are the key new concepts involved in the K_4 case.

Definition 3.1. (i) For a cube $C \in \mathcal{C}_1$, we define its 1-codimension as $d_1(C) = d(C)$. For a cube $C \in \mathcal{C}_2$ contained in a cube $C_1 \in \mathcal{C}_1$, we define its 1-codimension as $d_1(C) = d(C_1)$ and its 2-codimension as $d_2(C) = d(C) - d(C_1)$.

- (ii) Suppose $C, C' \in \mathcal{C}$ are adjacent. We say that they have level 1 adjacency if the level 1 cubes containing C and C' are different. Otherwise, they have level 2 adjacency (note that C and C' are both level 2 cubes in the second case).

After finding the double tiling, in order to find an embedding of the cube, we only consider the level 2 tiling (as explained above, we need to go through a level 1 tiling in order to obtain a ‘nice’ level 2 tiling). For two adjacent cubes $C, C' \in \mathcal{C}_2$, the definition of the corresponding sets $S_C, S_{C'}$ being in ‘good relation’ (see the definition of proper assignment in Section 3.2) now depends on the type of adjacency they have. If they have level 1 adjacency, then the good relation will be defined in terms of their 1-codimensions and if they have level 2 adjacency, then it will be in terms of their 2-codimensions. By doing this, we can overcome the above mentioned difficulty of not having enough control on the blue edges between S_C and $S_{C'}$. Afterwards, we proceed as in the previous section, by imposing a maximum degree condition between cubes in \mathcal{C}_2 and then embedding the cube.

We now provide the details of the proof. Our argument works for $n \geq 32$. Consider a coloring of the edges of the complete graph K_N on the vertex set $[N]$ for $N \geq 2^{46} \cdot 2^n$ with two colors, red and blue, and assume that there is no blue K_4 . We prove that this coloring contains a red Q_n .

3.1 Preprocessing the coloring

Note that we have $n \geq 2 \log n + 21$ for $n \geq 32$. We use the following procedure to construct our first level \mathcal{S}_1 of subsets of $[N]$:

For each $d = 0, 1, 2, \dots, \log n + 18$, if there exists a set S which induces a blue triangle-free graph of order exactly $2^{18} \cdot 2^{n-d}$, then arbitrarily choose one, add it to the family \mathcal{S}_1 , and remove the vertices of S from $[N]$. We define the *1-codimension* $d_1(S)$ of such a set as $d_1(S) = d$. When there are no more such sets, continue to the next value of d . If, after running through all values of d , $\sum_{S \in \mathcal{S}} |S| < \frac{N}{2}$, then add the set of remaining vertices to \mathcal{S} and declare it to be an *exceptional set* with 1-codimension zero (note that this set has size at least $\frac{N}{2} \geq 2^{18} \cdot 2^n$).

Now we perform a similar decomposition for each set in \mathcal{S}_1 to construct our second level \mathcal{S}_2 . For each set $S_1 \in \mathcal{S}_1$, consider the following procedure (suppose that S_1 has codimension $d_1 = d(S_1)$):

1. If S_1 is not exceptional: for each $d = 0, 1, 2, \dots, \log n + 3$, if there exists a subset $S \subset S_1$ which induces a red clique of order exactly $8 \cdot 2^{n-d_1-d}$, then arbitrarily choose one, add it to the family \mathcal{S}_2 , and remove the vertices of the clique from S_1 . We define the *2-codimension* of such a set as $d_2(S) = d$ and the *1-codimension* as $d_1(S) = d_1$. When there are no more such red cliques, continue to the next value of d . If, after running through all values of d , $\sum_{S \in \mathcal{S}_2, S \subset S_1} |S| < \frac{|S_1|}{2}$, then add the set of remaining vertices to \mathcal{S}_2 , and declare it to be an *exceptional set* with 2-codimension zero and 1-codimension d_1 (note that this set has size at least $\frac{|S_1|}{2} \geq 8 \cdot 2^{n-d_1}$).
2. If S_1 is exceptional: add S_1 to the family \mathcal{S}_2 . Define the 2-codimension of $S_1 \in \mathcal{S}_2$ as zero and its 1-codimension as zero.

Let $\mathcal{S} = \mathcal{S}_1 \cup \mathcal{S}_2$ (we suppose that \mathcal{S} is a multi-family and, thus, if a set is in both \mathcal{S}_1 and \mathcal{S}_2 , then we have two copies of this set in \mathcal{S} , however they can be distinguished by their levels). For a set $S \in \mathcal{S}$, define its *codimension* as $d(S) = \sum_{i \leq \ell(S)} d_i(S)$.

- Proposition 3.1.** (i) $\sum_{S \in \mathcal{S}_1} |S| \geq \frac{N}{2}$ and, for every $S_1 \in \mathcal{S}_1$, we have $\sum_{S \in \mathcal{S}_2, S \subset S_1} |S| \geq \frac{|S_1|}{2}$.
(ii) For an integer $i \geq 1$, let $X = \bigcup_{S \in \mathcal{S}_1, d_1(S) \geq i} S$. Then each vertex $v \in [N]$ has at most $2^{19} \cdot 2^{n-i}$ blue neighbors in X .
(iii) For a set $S_1 \in \mathcal{S}_1$ and an integer $i \geq 1$, let

$$X = \bigcup_{S \in \mathcal{S}_2, S \subset S_1, d_2(S) \geq i} S.$$

Then each vertex $v \in S_1$ has at most $16 \cdot 2^{n-d_1(S_1)-i}$ blue neighbors in X .

- (iv) For every set $S \in \mathcal{S}_2$, the subgraph induced by S has maximum blue degree at most $\frac{2^{n-d(S)}}{n}$.

Proof. Part (i) is clear and we omit the proof of (ii), since its proof is similar to that of part (i) of Proposition 2.1.

(iii) Suppose that S_1 is an exceptional set. Then S_1 is the unique set S satisfying $S \in \mathcal{S}_2$ and $S \subset S_1$. Since $d_2(S_1) = 0$, X is an empty set for every $i \geq 1$. Thus the conclusion follows.

Now suppose that S_1 is not an exceptional set. Since S_1 induces a blue triangle-free set, a blue neighborhood in S_1 of a vertex $v \in S_1$ forms a red clique. Thus, if a vertex $v \in S_1$ has more than $16 \cdot 2^{n-d_1(S_1)-i}$ blue neighbors in X , then inside the neighborhood of v we can find a set inducing a red clique of size at least $8 \cdot 2^{n-d_1(S_1)-(i-1)}$. However, this set must have been added in the previous step. Therefore, there are no such vertices.

(iv) Suppose that $S \subset S_1$ with $S_1 \in \mathcal{S}_1$. If both S and S_1 are not exceptional sets, then S induces a red clique and the conclusion follows. Suppose now that S is exceptional and S_1 is not. Then, as in (iii), we can see that there is no vertex in S which has at least $8 \cdot 2^{n-d_1(S)-(\log n+3)} = \frac{2^{n-d(S)}}{n}$ blue neighbors in S . One can similarly handle the case when S_1 is exceptional, since we have $S = S_1$ in this case. \square

3.2 Tiling the cube

The following proposition is similar to Proposition 2.2 (we omit its proof).

Proposition 3.2. Let $\mathcal{C} = \mathcal{C}_1 \cup \mathcal{C}_2$ be a double tiling, and let $C \in \mathcal{C}$ be a special cube of codimension d .

- (i) If C is a level 1 cube, then, for each $\ell = 1, 2$, there are at most $d_1(C)$ special cubes of level ℓ and codimension at most d which are adjacent to C (and C has level 1 adjacency with all these cubes).
(ii) If C is a level 2 cube, then there are at most $d_2(C)$ special cubes of codimension at most d which have level 2 adjacency with C (they are necessarily of level 2) and, for $\ell = 1, 2$, at most $d_1(C)$ cubes of level ℓ and codimension at most d which have level 1 adjacency with C .

Our double tiling \mathcal{C} of the cube Q_n will be constructed in correspondence with the family \mathcal{S} constructed in Section 3.1. We will construct \mathcal{C} by finding cubes of the tiling one by one. We slightly abuse notation and use \mathcal{C} also to denote the ‘partial’ double tiling, where only part of the cube Q_n is covered. Ideally, we would like to construct \mathcal{C} by constructing the level 1 tiling \mathcal{C}_1 first and then the level 2 tiling \mathcal{C}_2 . However, as we will soon see, it turns out that constructing the tiling in increasing order of codimension is more effective than in increasing order of level. At each step, we find a

subcube C which covers some non-covered part of Q_n and assign it to some set $S_C \in \mathcal{S}$. We say that such an assignment is *proper* if the following properties hold.

Proper assignment.

1. $\ell(C) = \ell(S_C)$ and, for $\ell = \ell(C)$, we have $d_\ell(C) = d_\ell(S_C)$.
2. If $C_2 \subset C_1$ with $C_1 \in \mathcal{C}_1$ and $C_2 \in \mathcal{C}_2$, then $S_{C_2} \subset S_{C_1}$.
3. Suppose that C is adjacent to some cube C' which is already in the tiling, where C' is assigned to $S_{C'}$ and C and C' have level ρ adjacency. Then the number of blue edges in the bipartite graph induced by S_C and $S_{C'}$ is at most $|S_C||S_{C'}| \cdot (8\delta)^{\ell(C)+\ell(C')-6}$, where $\delta = \max\{d_\rho(C), d_\rho(C')\}$.

Our algorithm for finding the tiling \mathcal{C} and the corresponding sets in \mathcal{S} is as follows.

Tiling Algorithm. At each step, consider all possible special cubes C which

- (a) can be added to \mathcal{C} to extend the partial tiling,
- (b) have $d(C) \geq d(C')$ for all $C' \in \mathcal{C}$, and
- (c) for which there exists a set $S \in \mathcal{S}$ which has not yet been assigned and such that assigning C to S gives a proper assignment.

Take a cube C_0 of minimum codimension satisfying (a), (b) and (c) and add it to the tiling. Assign C_0 to the set $S_{C_0} \in \mathcal{S}$ given by (c).

Condition (a) is equivalent to saying that either C is disjoint from all cubes in \mathcal{C}_1 or is contained in some cube in \mathcal{C}_1 and is disjoint from all cubes in \mathcal{C}_2 . The following proposition shows that the algorithm will terminate successfully.

Proposition 3.3. *If the tiling is not complete, then the algorithm always chooses a cube from a non-empty collection.*

Proof. Suppose that in the previous step we embedded some cube of codimension d (let $d = 0$ for the first iteration of the algorithm). Since the tiling is not complete, there exists a vertex $(a_1, \dots, a_n) \in Q_n$ which is not covered twice.

Case 1: (a_1, \dots, a_n) is not covered by any of the cubes in \mathcal{C}_1 .

Let \mathcal{S}'' be the subfamily of \mathcal{S}_1 consisting of sets which are already assigned to some cube in \mathcal{C}_1 , and let $\mathcal{S}' = \mathcal{S}_1 \setminus \mathcal{S}''$. If \mathcal{S}'' contains a set of codimension zero, then the corresponding cube is the whole cube Q_n , contradicting the fact that (a_1, \dots, a_n) is not covered. Thus we may assume that there is no set of codimension zero in \mathcal{S}'' , from which it follows that $|S_C| = 2^{18}|C|$ for all $C \in \mathcal{C}_1$. Note that

$$\left| \bigcup_{S \in \mathcal{S}'} S \right| = \left| \bigcup_{S \in \mathcal{S}_1} S \right| - \left| \bigcup_{S \in \mathcal{S}''} S \right| \geq \frac{N}{2} - \left| \bigcup_{C \in \mathcal{C}} S_C \right| \geq 2^{45} \cdot 2^n - 2^{18} \cdot 2^n > 2^{44} \cdot 2^n.$$

For each i , let $\mathcal{S}'_i = \{S \in \mathcal{S}' : d(S) = i\}$. Suppose that $|\mathcal{S}'_i| \leq 2^{14}i^5$ for all i . Then, since each set in \mathcal{S}'_i has size $2^{18} \cdot 2^{n-i}$, we have

$$\left| \bigcup_{S \in \mathcal{S}'} S \right| = \sum_i \sum_{S \in \mathcal{S}'_i} |S| \leq \sum_i 2^{14}i^5 \cdot 2^{18} \cdot 2^{n-i} = 2^{32} \cdot 2^n \sum_i \frac{i^5}{2^i} < 2^{44} \cdot 2^n,$$

which is a contradiction (we used the fact that $\sum_i \frac{i^5}{2^i} < 2^{12}$). Therefore, there exists an index i for which $|\mathcal{S}'_i| > 2^{14}i^5$. We have $i > 0$ since otherwise the set of codimension zero would have been added to the tiling in the first step.

Let $C = (a_1, \dots, a_i, *, \dots, *)$. For each $\ell = 1, 2$, let \mathcal{A}_ℓ be the family of cubes of level ℓ in our partial double tiling which are adjacent to C and have codimension at most i . By Proposition 3.2, we have $|\mathcal{A}_\ell| \leq i$ for each ℓ . For a cube $A \in \mathcal{A}_\ell$, we say that a set $S \in \mathcal{S}'_i$ is *bad* for A if there are at least $\frac{1}{(8i)^{5-\ell}}|S_A||S|$ blue edges between S_A and S . Otherwise, we say that S is *good* for A . We claim that there are at most $2^{13}i^4$ sets in \mathcal{S}'_i which are bad for each fixed A .

Let $X_i = \bigcup_{S \in \mathcal{S}'_i} S$ and note, by Proposition 3.1, that there are at most $|S_A| \cdot 2^{19} \cdot 2^{n-i}$ blue edges between the sets S_A and X_i . Each set $S \in \mathcal{S}'_i$ which is bad for A accounts for at least

$$\frac{1}{(8i)^{5-\ell}}|S_A||S| \geq \frac{2^{18} \cdot 2^{n-i}}{(8i)^4}|S_A| = \frac{2^6 \cdot 2^{n-i}}{i^4}|S_A|$$

such blue edges (note that $|S| = 2^{18} \cdot 2^{n-i}$). Therefore, in total, there are at most $2^{13} \cdot i^4$ sets in \mathcal{S}'_i which are in bad relation with A , as claimed above. Since $|\mathcal{A}_1 \cup \mathcal{A}_2| \leq 2i$ and $|\mathcal{S}'_i| > 2^{14}i^5$, there exists $S \in \mathcal{S}'_i$ which is good for all cubes in $\mathcal{A}_1 \cup \mathcal{A}_2$.

In order to show that C satisfies (a) and (b), it suffices to verify that $i \geq d$, since $i \geq d$ implies the fact that C is disjoint from all the other cubes in \mathcal{C}_1 (if C intersects some other cube, then that cube must contain C and therefore also contains (a_1, \dots, a_n) by Proposition 2.2). Furthermore, if this is the case, assigning C to S is a proper assignment (thus we have (c)).

Now suppose, for the sake of contradiction, that $i < d$ and consider the time t immediately after we last embedded a cube of codimension at most i . At time t , since (a_1, \dots, a_n) was not covered, the cubes in \mathcal{C} are disjoint from C . Moreover, the set of cubes of codimension at most i which are adjacent to C is the same as at current time. Therefore, C could have been added to the tiling at time t as well, and this contradicts the fact that we always choose a cube of minimum codimension. Thus we have $i \geq d$, as claimed.

Case 2: (a_1, \dots, a_n) is covered by a cube $C_1 \in \mathcal{C}_1$ but not by a cube in \mathcal{C}_2 .

In this case, in order to assign some cube $C \subset C_1$ containing (a_1, \dots, a_n) to a subset of S_{C_1} in \mathcal{S}_2 , there are two types of adjacency that one needs to consider. Let $C_2 = (a_1, \dots, a_d, *, \dots, *)$ and temporarily consider it as a level 2 cube. We first consider the cubes which are 1-adjacent to C_2 and remove all the subsets of S_{C_1} which are ‘bad’ for these cubes. The reason we consider cubes which are adjacent to C_2 instead of C is because we do not know what C will be at this point. Note that, depending on the choice of C , some of the cubes that are 1-adjacent to C_2 may not be 1-adjacent to C . However, since we choose $C \subset C_2$ (as we will show later in the proof), the set of cubes which are 1-adjacent to C will be a subset of the set of cubes which are 1-adjacent to C_2 . Moreover, if C' is a cube 1-adjacent to both C and C_2 , then $\delta(C', C) = \delta(C', C_2)$. Having removed all ‘bad’ subsets, we have enough information to determine C . Then, by considering the relation of the remaining subsets of S_{C_1} to cubes which are 2-adjacent to C , we can find a set S_C that can be assigned to C . Note that unlike in the previous case, even if C' is a cube 2-adjacent to both C and C_2 , we do not necessarily have $\delta(C', C) = \delta(C', C_2)$. For this reason, it turns out to be crucial that we have already determined C before the second step.

Let $d_1 = d(C_1)$ and let $\mathcal{A}^{(1)}$ be the subfamily of our partial embedding \mathcal{C} consisting of cubes which are 1-adjacent to C_2 . Note that, by Proposition 3.2(ii), there are at most d_1 cubes which are 1-adjacent

to C_2 in each level and thus $|\mathcal{A}^{(1)}| \leq 2d_1$. Let $\mathcal{F} = \{S \in \mathcal{S}_2 : S \subset S_{C_1}\}$. By Proposition 3.1, we have $\left| \bigcup_{S \in \mathcal{F}} S \right| \geq \frac{|S_{C_1}|}{2}$. We say that a set $S \in \mathcal{F}$ is *bad* for a cube $A \in \mathcal{A}^{(1)}$ if there are at least $(8\delta_A)^{\ell(S)+\ell(S_A)-6}|S||S_A| = \frac{|S||S_A|}{(8\delta_A)^{4-\ell(A)}}$ blue edges between S and S_A (where $\delta_A = \max\{d_1, d_1(A)\}$). Otherwise, we say that S is *good* for A . For each fixed A , let \mathcal{F}_A be the subfamily of \mathcal{F} consisting of sets which are bad for A . By the properness of the assignment up to this point, we know that there are at most $\frac{|S_{C_1}||S_A|}{(8\delta_A)^{5-\ell(A)}}$ blue edges between S_{C_1} and S_A for every $A \in \mathcal{A}^{(1)}$. Therefore, by counting the number of blue edges between S_{C_1} and S_A in two ways, we see that

$$\sum_{S \in \mathcal{F}_A} \frac{|S||S_A|}{(8\delta_A)^{4-\ell(A)}} \leq \frac{|S_{C_1}||S_A|}{(8\delta_A)^{5-\ell(A)}},$$

from which we have $\sum_{S \in \mathcal{F}_A} |S| \leq \frac{1}{8\delta_A} |S_{C_1}| \leq \frac{1}{8d_1} |S_{C_1}|$.

Let \mathcal{F}'' be the subfamily of \mathcal{F} of sets which are already assigned to some cube in \mathcal{C}_2 . There are no sets of relative codimension zero in \mathcal{F}'' since this would imply that (a_1, \dots, a_n) is covered by a cube in \mathcal{C}_2 . It thus follows that for every $C \in \mathcal{C}_2$ such that $C \subset C_1$, we have $|S_C| = 8|C|$ and

$$\left| \bigcup_{S \in \mathcal{F}''} S \right| = \sum_{S \in \mathcal{F}''} |S| \leq 8 \sum_{C \subset C_1} |C| \leq 8|C_1| = 8 \cdot 2^{n-d_1}.$$

Let $\mathcal{S}' = \mathcal{F} \setminus (\mathcal{F}'' \cup \bigcup_{A \in \mathcal{A}^{(1)}} \mathcal{F}_A)$ be the subfamily of $\mathcal{F} = \{S \in \mathcal{S}_2 : S \subset S_{C_1}\}$ of sets which are not assigned to any cubes yet and are good for all the cubes in $\mathcal{A}^{(1)}$. Since $|S_{C_1}| \geq 2^{18}|C_1| = 2^{18} \cdot 2^{n-d_1}$, we have

$$\begin{aligned} \left| \bigcup_{S \in \mathcal{S}'} S \right| &\geq \left| \bigcup_{S \in \mathcal{F}} S \right| - \left| \bigcup_{S \in \mathcal{F}''} S \right| - \sum_{A \in \mathcal{A}^{(1)}} \left| \bigcup_{S \in \mathcal{F}_A} S \right| \\ &\geq \frac{|S_{C_1}|}{2} - 8 \cdot 2^{n-d_1} - 2d_1 \cdot \frac{|S_{C_1}|}{8d_1} > 2^{15} \cdot 2^{n-d_1}. \end{aligned}$$

For each $i \geq 0$, let $\mathcal{S}'_i = \{S \in \mathcal{S}' : d_2(S) = i\} = \{S \in \mathcal{S}' : d(S) = d_1 + i\}$. Suppose that $|\mathcal{S}'_i| \leq 128i^3$ for all i . Then, since we have $|S| = 8 \cdot 2^{n-d_1-i}$ for all $S \in \mathcal{S}'_i$,

$$\left| \bigcup_{S \in \mathcal{S}'} S \right| = \sum_i \sum_{S \in \mathcal{S}'_i} |S| \leq \sum_i 128i^3 \cdot 8 \cdot 2^{n-d_1-i} < 2^{15} \cdot 2^{n-d_1},$$

which is a contradiction. Therefore, there exists an index i for which $|\mathcal{S}'_i| > 128 \cdot i^3$.

Let $C = (a_1, \dots, a_{d_1+i}, *, \dots, *)$ and consider it as a level 2 cube. By Proposition 3.2(ii), there are at most i cubes in \mathcal{C} which have level 2 adjacency with C and have codimension at most $d_1 + i$. Let $\mathcal{A}^{(2)}$ be the family consisting of these cubes. For a cube $A \in \mathcal{A}^{(2)}$, we say that a set $S \in \mathcal{S}'_i$ is *bad* for A if there are at least $(8i)^{\ell(S_A)+\ell(S)-6}|S_A||S| = \frac{1}{(8i)^2}|S_A||S|$ blue edges between S_A and S . Otherwise, we say that S is *good* for A . We claim that there are at most $128i^2$ sets in \mathcal{S}'_i which are bad for each fixed A .

Let $X_i = \bigcup_{S \in \mathcal{S}'_i} S$ and note, by Proposition 3.1, that there are at most $|S_A| \cdot 16 \cdot 2^{n-i}$ blue edges between the sets S_A and X_i . Each set $S \in \mathcal{S}'_i$ which is bad for A accounts for at least

$$\frac{1}{(8i)^2} |S_A||S| = \frac{8 \cdot 2^{n-i}}{64i^2} |S_A| = \frac{2^{n-i}}{8i^2} |S_A|$$

such blue edges (note that $|S| = 8 \cdot 2^{n-i}$). Therefore, in total, there are at most $128i^2$ sets in \mathcal{S}'_i which are bad for A , as claimed above. Since there are at most i sets in $\mathcal{A}^{(2)}$ and $|\mathcal{S}'_i| > 128i^3$, there exists a set $S \in \mathcal{S}'_i$ which is good for all the cubes $A \in \mathcal{A}^{(2)}$.

In order to show that C satisfies (a) and (b), it suffices to verify that $d_1 + i \geq d$, since this implies the fact that C is disjoint from all the other cubes in \mathcal{C}_2 (note that $C \subset C_1$ and if C intersects some other cube of level 2, then that cube must contain C and therefore also contains (a_1, \dots, a_n) by Proposition 2.2). Furthermore, if this is the case, assigning C to S is a proper assignment (thus we have (c)).

Now suppose, for the sake of contradiction, that $d_1 + i < d$ and consider the time t immediately after we last embedded a cube of codimension at most $d_1 + i$. At time t , since $d_1 + i \geq d_1 = d(C_1)$, the cube C_1 was already embedded and, since there are no cubes of level 2 covering (a_1, \dots, a_n) , the cubes in \mathcal{C}_2 are disjoint from C . A cube in the partial embedding at time t , which is of codimension at most $d_1 + i$, is 1-adjacent to C if and only if it is 1-adjacent to C_2 . Hence, the family of cubes which are adjacent to C at time t is a subfamily of $\mathcal{A}^{(1)} \cup \mathcal{A}^{(2)}$. Therefore, C could have been added to the tiling at time t as well, contradicting the fact that we always choose a cube of minimum codimension. Thus we have $d_1 + i \geq d$, as claimed. \square

Note that as an outcome of our algorithm, we obtain a tiling \mathcal{C} such that for every pair of adjacent cubes $C, C' \in \mathcal{C}$, we have control on the number of blue edges between S_C and $S_{C'}$ (as given in the definition of proper assignment).

3.3 Imposing a maximum degree condition

As in the triangle case, we now impose certain maximum degree conditions between the sets S_C for $C \in \mathcal{C}$. It suffices to impose maximum degree conditions between sets of level 2.

For a set $C \in \mathcal{C}_2$ of codimension $d = d(C)$, and relative codimensions $d_1 = d_1(C), d_2 = d_2(C)$, recall that we have a set $S_C \in \mathcal{S}$ such that $|S_C| \geq 8 \cdot 2^{n-d}$. Let \mathcal{A}_ρ be the family of cubes in \mathcal{C}_2 with codimension at most d which have level ρ adjacency with C . For each $A \in \mathcal{A}_\rho$, let $\delta_{A,\rho} = \max\{d_\rho, d_\rho(A)\}$. By Proposition 3.2(ii), we have $|\mathcal{A}_\rho| \leq d_\rho$ for each $\rho = 1, 2$. For each $A \in \mathcal{A}_\rho$, there are at most $\frac{1}{64\delta_{A,\rho}^2} |S_C| |S_A|$ blue edges between S_C and S_A .

Now for $\rho = 1, 2$ and each $A \in \mathcal{A}_\rho$, remove all the vertices in S_C which have at least $\frac{1}{8\delta_{A,\rho}} |S_A|$ blue neighbors in S_A and let T_C be the subset of S_C left after these removals. Since there are at most $\frac{1}{64\delta_{A,\rho}^2} |S_C| |S_A|$ blue edges between S_C and S_A , we remove at most $\frac{|S_C|}{8\delta_{A,\rho}}$ vertices from S_C for each set $A \in \mathcal{A}_\rho$. Thus the resulting set T_C is of size at least

$$|T_C| \geq |S_C| - d_1 \cdot \frac{|S_C|}{8\delta_{A,1}} - d_2 \cdot \frac{|S_C|}{8\delta_{A,2}} \geq \frac{|S_C|}{2} \geq 4 \cdot 2^{n-d}.$$

For each $A \in \mathcal{A}_\rho$, all the vertices in T_C have blue degree at most $\frac{1}{8\delta_{A,\rho}} |S_A| \leq \frac{1}{4\delta_{A,\rho}} |T_A|$ in the set T_A . Thus we obtain the following property.

Maximum degree condition. Let $C, C' \in \mathcal{C}_2$ be a pair of cubes having level ρ adjacency with $d(C) \geq d(C')$. Then every vertex in T_C has at most $\frac{1}{4\delta_\rho} |T_{C'}|$ blue neighbors in the set $T_{C'}$ (where $\delta_\rho = \max\{d_\rho(C), d_\rho(C')\}$).

3.4 Embedding the cube

We now show how to embed Q_n . Recall that we found a double tiling \mathcal{C} of Q_n . We will greedily embed the cubes in the level 2 tiling \mathcal{C}_2 one by one into their assigned sets from the family \mathcal{S}_2 , in decreasing order of their codimensions. If there are several cubes of the same codimension, then we arbitrarily choose the order between them.

Suppose that we are about to embed the cube $C \in \mathcal{C}_2$. Let $d = d(C)$, $d_1 = d_1(C)$ and $d_2 = d_2(C)$. We will greedily embed the vertices of C into the set $T_C \subseteq S_C$. Suppose that we are about to embed $x \in C$ and let $f : Q_n \rightarrow [N]$ denote the partial embedding of the cube Q_n obtained so far. For $\rho = 1, 2$, let A_ρ be the set of neighbors of x which are already embedded and belong to a cube other than C that has level ρ adjacency with C . Note that we have $|A_\rho| \leq d_\rho$ for $\rho = 1$ and 2 . Since we have so far only embedded cubes of codimension at least d , we have, for each $\rho = 1, 2$, that the vertices $f(v)$ for $v \in A_\rho$ have blue degree at most $\frac{1}{4d_\rho}|T_C|$ in the set T_C , by the maximum degree condition imposed in Section 3.3. Together, these neighbors forbid at most

$$d_1 \cdot \frac{1}{4d_1}|T_C| + d_2 \cdot \frac{1}{4d_2}|T_C| \leq \frac{1}{2}|T_C|.$$

vertices of T_C from being the image of x .

In addition, x has at most $n - d$ neighbors which are already embedded and belong to C . By Proposition 3.1, for each such vertex v , $f(v)$ has blue degree at most $\frac{2^{n-d}}{n}$ in the set T_C . Together, these neighbors forbid at most 2^{n-d} vertices of T_C from being the image of x . Finally, there are at most $2^{n-d} - 1$ vertices in T_C which are images of some other vertex of C that is already embedded. Therefore, the number of vertices in T_C into which we cannot embed x is at most

$$\frac{1}{2}|T_C| + 2^{n-d} + (2^{n-d} - 1) < |T_C|,$$

where the inequality follows since $|T_C| \geq 4 \cdot 2^{n-d}$. Hence, there exists a vertex in T_C which we can choose as an image of x to extend the current partial embedding of the cube. Repeating this procedure until we embed the whole cube Q_n completes the proof.

4 General case

In this section, we further extend the arguments presented so far to prove the main theorem. The framework is very similar to that of the previous sections, where we find multiple levels of tiling of Q_n . We begin by preprocessing the coloring to find families of sets $\mathcal{S} = \mathcal{S}_0 \cup \mathcal{S}_1 \cup \dots \cup \mathcal{S}_{s-2}$ such that, for $\ell \geq 1$, sets in \mathcal{S}_ℓ are subsets of sets in $\mathcal{S}_{\ell-1}$ which do not contain blue $K_{s-\ell}$ (except for some special cases). We refer to the sets in \mathcal{S}_ℓ as *level ℓ sets* and, for a set $S \in \mathcal{S}_\ell$, we let its *level* be $\ell(S) = \ell$. Note that, unlike in the cases where $s = 3, 4$, we have an additional level, level zero. Level zero will consist of a single set $[N]$ and it is there merely for technical reasons.

We then seek a corresponding (multiple level) tiling \mathcal{C} of the cube Q_n . Define an $(s-1)$ -fold tiling (or $(s-1)$ -tiling, for short) of Q_n as a collection of $s-1$ tilings $\mathcal{C}_0 \cup \mathcal{C}_1 \cup \dots \cup \mathcal{C}_{s-2}$, where \mathcal{C}_0 is the trivial tiling consisting of the unique cube Q_n , and, for all $\ell \geq 1$, \mathcal{C}_ℓ is a refined tiling of $\mathcal{C}_{\ell-1}$ (i.e. for every $C \in \mathcal{C}_\ell$, there exists $C' \in \mathcal{C}_{\ell-1}$ such that $C \subset C'$). We refer to cubes in \mathcal{C}_ℓ as *level ℓ cubes* and, for a cube $C \in \mathcal{C}_\ell$, define its *level* as $\ell(C) = \ell$. We will construct the tiling by finding cubes $C \in \mathcal{C}$

and assigning each of them to some set $S_C \in \mathcal{S}$. Informally, this means that the subcube C of Q_n will be found in the S_C part of our graph. Note that the trivial level zero cube Q_n gets assigned to the trivial level zero set $[N]$ and this fits the heuristic.

For the rest of this section, we assume that $s \geq 5$. Let $c = s^{15s}$ and suppose that $N \geq c^s 2^n = s^{15s^2} 2^n$. We will later use the following estimate.

Lemma 4.1. *For every positive integer s , $\sum_{i=1}^{\infty} \frac{i^s}{2^i} \leq 2s^s$.*

Proof. Let $(x)_t = x(x-1)\dots(x-t+1)$, $X_t = \sum_{i=1}^{\infty} \frac{i^t}{2^i}$ and $Y_t = \sum_{i=1}^{\infty} \frac{(i)_t}{2^i}$ for non-negative integers t . The *Stirling number $S(t, k)$ of the second kind* is the number of ways to partition a set of t objects into k non-empty subsets. These numbers satisfy the following well-known identity $x^t = \sum_{k=0}^t S(t, k)(x)_k$ (see, e.g., [34], Chapter 1.4). This implies the identity

$$X_t = \sum_{k=0}^t S(t, k) Y_k.$$

By taking the derivative k times of both sides of the equality $(1-z)^{-1} = \sum_{i \geq 0} z^i$, note that

$$k! \cdot (1-z)^{-(k+1)} = \sum_{i \geq 1} i(i-1)\dots(i-k+1)z^{i-k}.$$

By multiplying both sides by z^k and substituting $z = 1/2$ we have that $Y_k = 2k!$. This, together with the above identity, implies that $X_t = \sum_{k=0}^t 2k! S(t, k)$. Let T_s be the number of partitions of a set of s objects into labelled nonempty subsets. By counting over the size k of the partition, we have the identity $T_s = \sum_{k=0}^s k! S(s, k) = X_s/2$. As each partition counted by T_s is determined by the vector of labels of the sets containing each object, and there are at most s such labels for each partition, we have $T_s \leq s^s$ and the desired inequality follows.

(Although it will not be needed, we remark that there is an explicit formula

$$X_t = 2 \sum_{k=0}^t \sum_{j=0}^k (-1)^{k-j} \binom{k}{j} j^t$$

which follows from substituting in the well-known identity $S(t, k) = \frac{1}{k!} \sum_{j=0}^k (-1)^{k-j} \binom{k}{j} j^t$. \square

4.1 Preprocessing the coloring

Let $\mathcal{S}_0 = \{[N]\}$ and $[N]$ be the unique set of level zero and codimension zero (denoted as $\ell(S) = 0$ and $d(S) = 0$). We construct the levels one at a time. Once we finish constructing $\mathcal{S}_{\ell-1}$, for each set $S' \in \mathcal{S}_{\ell-1}$, we use the following procedure to construct sets belonging to the ℓ -th level \mathcal{S}_{ℓ} :

1. If S' is not exceptional: for each $d = 0, 1, 2, \dots, \log n + s \log c$, if there exists a set $S \subset S'$ which induces a blue $K_{s-\ell}$ -free graph of order exactly $c^{s-\ell} \cdot 2^{n-d(S')-d}$, then arbitrarily choose one, add it to the family \mathcal{S}_{ℓ} and remove the vertices of S from S' . We define the ℓ -codimension $d_{\ell}(S)$ of such a set as $d_{\ell}(S) = d$ and, for $i = 1, \dots, \ell - 1$, we define the i -codimension $d_i(S)$ as $d_i(S) = d_i(S')$. Let the *codimension* of S be $d(S) = \sum_{i=1}^{\ell} d_i(S)$.

When there are no more such sets, continue to the next value of d . If, after running through all values of d , $\sum_{S \in \mathcal{S}_\ell, S \subset S'} |S| < \frac{|S'|}{2}$, then add the set of remaining vertices to \mathcal{S}_ℓ and declare it to be an *exceptional set* with ℓ -codimension zero (note that this set has size at least $\frac{|S'|}{2} = \frac{c^{s-(\ell-1)} \cdot 2^{n-d(S')}}{2} \geq c^{s-\ell} \cdot 2^{n-d(S')}$).

2. If S' is exceptional: add S' to the family \mathcal{S}_ℓ and define its ℓ -codimension as zero.

Let $\mathcal{S} = \bigcup_{\ell=0}^{s-2} \mathcal{S}_\ell$ (we suppose that \mathcal{S} is a multi-family and if a set appears multiple times we distinguish them by their levels). The following proposition is similar to Proposition 3.1. We omit its proof.

Proposition 4.1. (i) For $1 \leq \ell \leq s-2$ and a set $S' \in \mathcal{S}_{\ell-1}$, $\sum_{S \in \mathcal{S}_\ell, S \subset S'} |S| \geq \frac{|S'|}{2}$.
(ii) For integers $1 \leq \ell \leq s-2$ and $i \geq 1$, let S' be a set of level $\ell-1$ and let

$$X = \bigcup_{S \in \mathcal{S}_\ell, S \subset S', d_\ell(S) \geq i} S.$$

Then each vertex $v \in S'$ has at most $2c^{s-\ell} \cdot 2^{n-d(S')-i}$ blue neighbors in X .

(iii) For every set $S \in \mathcal{S}_{s-2}$, the subgraph induced by S has maximum blue degree at most $\frac{2^{n-d(S)}}{n}$.

4.2 Tiling the cube

In this subsection, we find an $(s-1)$ -tiling of Q_n . Recall that in the previous section, we had to control the blue edge densities between adjacent cubes in the tiling. The parameter that governed the control of these densities was defined in terms of the ρ -codimension, where ρ was the level of adjacency of these cubes. Below we generalize this concept.

Definition 4.1. Let \mathcal{C} be an $(s-1)$ -tiling and let $C, C' \in \mathcal{C}$ be two adjacent cubes.

- (i) The level of adjacency $\rho(C, C')$ is the minimum ℓ such that the cubes of level ℓ containing C and C' are distinct. We say that C and C' are ρ -adjacent if $\rho(C, C') = \rho$.
- (ii) The dominating parameter $\delta(C, C')$ is $\max\{d_\rho(C), d_\rho(C')\}$, where $\rho = \rho(C, C')$.

Note that the level of adjacency of two cubes C and C' is at most $\min\{\ell(C), \ell(C')\}$. The following proposition is similar to Proposition 3.2 and we omit its proof.

Proposition 4.2. Let \mathcal{C} be an $(s-1)$ -tiling and let $C \in \mathcal{C}$ be a level ℓ special cube of codimension d . For each $\rho = 1, 2, \dots, \ell$ and $\ell' = 1, 2, \dots, s-2$, the number of special cubes of level ℓ' and codimension at most d which are ρ -adjacent to C is at most $d_\rho(C)$.

Our $(s-1)$ -tiling \mathcal{C} of the cube Q_n will be constructed in correspondence with the family \mathcal{S} constructed in Section 4.1. We will construct \mathcal{C} by finding cubes of the tiling one by one. We slightly abuse notation and use \mathcal{C} also to denote the ‘partial’ $(s-1)$ -tiling, where only part of the cube Q_n is covered. At each step, we find a subcube C which covers some non-covered part of Q_n and assign it to some set $S_C \in \mathcal{S}$. We say that such an assignment is *proper* if the following properties hold.

Proper assignment.

1. $\ell(C) = \ell(S_C)$ and, for $\ell = \ell(C)$, we have $d_\ell(C) = d_\ell(S_C)$.

2. If $C \subset C'$ for $C' \in \mathcal{C}$, then $S_C \subset S_{C'}$.
3. Suppose that C is adjacent to some cube C' already in the tiling, where C' is assigned to $S_{C'}$, and that C and C' have level ρ adjacency. Then the number of blue edges in the bipartite graph induced by S_C and $S_{C'}$ is at most $|S_C||S_{C'}| \cdot (4s^2\delta(C, C'))^{\ell(C)+\ell(C')-2s}$.

We use S_C to denote the set in \mathcal{S} to which C is assigned. Our algorithm for finding the tiling \mathcal{C} and the corresponding sets in \mathcal{S} is as follows.

Tiling Algorithm. At each step, consider all possible cubes C which

- (a) can be added to \mathcal{C} to extend the partial tiling,
- (b) have $d(C) \geq d(C')$ for all $C' \in \mathcal{C}$, and
- (c) for which there exists a set $S \in \mathcal{S}$ which has not yet been assigned and such that assigning C to S gives a proper assignment.

Take a cube C_0 of minimum codimension satisfying (a), (b) and (c) and add it to the tiling. Assign C_0 to the set $S_{C_0} \in \mathcal{S}$ given by (c).

The following proposition shows that the algorithm will terminate successfully.

Proposition 4.3. *If the tiling is not complete, then the algorithm always chooses a cube from a non-empty collection.*

Proof. In the beginning, the algorithm will take Q_n as the level 0 tiling and will assign it to $[N]$, which is the unique set of level 0. Now suppose that in the previous step we embedded some cube of codimension d . Since the tiling is not complete, there exists a vertex $(a_1, \dots, a_n) \in Q_n$ which is not covered $s - 1$ times. Suppose that this vertex is covered ℓ times for $\ell \leq s - 2$ and let $C_0, \dots, C_{\ell-1}$ be the cubes of each level that cover it (note that $C_0 \supseteq C_1 \supseteq \dots \supseteq C_{\ell-1}$). Let $C_\ell = (a_1, \dots, a_d, *, \dots, *)$. We temporarily consider C_ℓ as a level ℓ cube.

In this case, in order to assign some cube $C \subset C_{\ell-1}$ containing (a_1, \dots, a_n) to a subset of $S_{C_{\ell-1}}$ in \mathcal{S}_ℓ , we first consider the cubes which are ρ -adjacent to C_ℓ for $\rho \leq \ell - 1$ and remove all the subsets of $S_{C_{\ell-1}}$ which are ‘bad’ for these cubes. The reason we consider cubes which are adjacent to C_ℓ instead of C is because we do not know what C will be at this point. Note that, depending on the choice of C , some of the cubes that are ρ -adjacent to C_ℓ may not be ρ -adjacent to C . However, since we choose $C \subset C_\ell$ (as we will show later in the proof), the set of cubes which are ρ -adjacent to C will be a subset of the set of cubes which are ρ -adjacent to C_ℓ . Moreover, for $\rho \leq \ell - 1$, if C' is a cube ρ -adjacent to both C and C_ℓ , then $\delta(C', C) = \delta(C', C_\ell)$. Having removed these ‘bad’ subsets, we have enough information to determine C . Then, by considering the relation of the remaining subsets of $S_{C_{\ell-1}}$ to cubes which are ℓ -adjacent to C , we can find a set S_C that can be assigned to C . Note that unlike in the previous case, for $\rho \geq \ell$, even if C' is a cube ρ -adjacent to both C and C_ℓ , we do not necessarily have $\delta(C', C) = \delta(C', C_\ell)$. For this reason, it turns out to be crucial that we have already determined C before the second step.

For each $\rho \leq \ell - 1$ and i , let $\mathcal{A}_i^{(\rho)}$ be the family of level i cubes from our partial embedding \mathcal{C} consisting of cubes which are ρ -adjacent to C_ℓ . Let $\mathcal{A}^{(\rho)} = \bigcup_i \mathcal{A}_i^{(\rho)}$. By Proposition 4.2, we have $|\mathcal{A}_i^{(\rho)}| \leq d_\rho(C_{\ell-1})$ for each i and thus $|\mathcal{A}^{(\rho)}| \leq s \cdot d_\rho(C_{\ell-1})$. Let $\mathcal{F} = \{S \in \mathcal{S}_\ell : S \subset S_{C_{\ell-1}}\}$. By

Proposition 4.1, we have $\left| \bigcup_{S \in \mathcal{F}} S \right| \geq \frac{|S_{C_{\ell-1}}|}{2}$. We say that a set $S \in \mathcal{F}$ is *bad* for a cube $A \in \mathcal{A}_i^{(\rho)}$ if there are at least $(4s^2\delta_A)^{\ell(S)+\ell(S_A)-2s}|S||S_A| = (4s^2\delta_A)^{\ell+i-2s}|S||S_A|$ blue edges between S and S_A (where $\delta_A = \delta(C_{\ell-1}, A)$ is the dominating parameter of $C_{\ell-1}$ and A). Otherwise, we say that S is *good* for A . For each fixed A , let \mathcal{F}_A be the subfamily of \mathcal{F} consisting of sets which are bad for A . By the properness of the assignment up to this point, we know that there are at most $(4s^2\delta_A)^{\ell+i-1-2s}|S_{C_{\ell-1}}||S_A|$ blue edges between $S_{C_{\ell-1}}$ and S_A for every $A \in \mathcal{A}_i^{(\rho)}$. Therefore, by counting the number of blue edges between $S_{C_{\ell-1}}$ and S_A in two ways, we see that

$$\sum_{S \in \mathcal{F}_A} (4s^2\delta_A)^{\ell+i-2s}|S||S_A| \leq (4s^2\delta_A)^{\ell+i-1-2s}|S_{C_{\ell-1}}||S_A|,$$

from which we have $\sum_{S \in \mathcal{F}_A} |S| \leq \frac{1}{4s^2\delta_A}|S_{C_{\ell-1}}| \leq \frac{1}{4s^2d_\rho(C_{\ell-1})}|S_{C_{\ell-1}}|$.

Let \mathcal{F}'' be the subfamily of \mathcal{F} of sets which are already assigned to some cube in \mathcal{C}_ℓ . There can be no sets of relative codimension zero in \mathcal{F}'' since this would imply that (a_1, \dots, a_n) is covered by a cube in \mathcal{C}_ℓ . It thus follows that for every $C \in \mathcal{C}_\ell$ such that $C \subset C_{\ell-1}$, we have $|S_C| = c^{s-\ell}|C|$ and

$$\left| \bigcup_{S \in \mathcal{F}''} S \right| = \sum_{S \in \mathcal{F}''} |S| \leq c^{s-\ell} \cdot |C_{\ell-1}|.$$

Let $\mathcal{S}' = \mathcal{F} \setminus (\mathcal{F}'' \cup \bigcup_\rho \bigcup_{A \in \mathcal{A}^{(\rho)}} \mathcal{F}_A)$ be the subfamily of \mathcal{F} of sets which are not assigned to any cubes yet and are good for all the cubes in $\mathcal{A}^{(\rho)}$ for all $\rho \leq \ell - 1$. Since $|S_{C_{\ell-1}}| \geq c^{s-\ell+1}|C_{\ell-1}|$, we have

$$\begin{aligned} \left| \bigcup_{S \in \mathcal{S}'} S \right| &\geq \left| \bigcup_{S \in \mathcal{F}} S \right| - \left| \bigcup_{S \in \mathcal{F}''} S \right| - \sum_\rho \sum_{A \in \mathcal{A}^{(\rho)}} \left| \bigcup_{S \in \mathcal{F}_A} S \right| \\ &\geq \frac{|S_{C_{\ell-1}}|}{2} - c^{s-\ell} \cdot |C_{\ell-1}| - \sum_\rho s d_\rho(C_{\ell-1}) \cdot \frac{|S_{C_{\ell-1}}|}{4s^2 d_\rho(C_{\ell-1})} > \frac{c^{s-\ell+1}}{5} \cdot |C_{\ell-1}|. \end{aligned}$$

For each $i \geq 0$, let $\mathcal{S}'_i = \{S \in \mathcal{S}' : d_\ell(S) = i\} = \{S \in \mathcal{S}' : d(S) = d(C_{\ell-1}) + i\}$. Suppose that $|\mathcal{S}'_i| \leq (4s^2i)^{2s+1}$ for all i . Then, since we have $|S| = c^{s-\ell}2^{n-d(C_{\ell-1})-i} = c^{s-\ell} \cdot 2^{-i}|C_{\ell-1}|$ for all $S \in \mathcal{S}'_i$,

$$\left| \bigcup_{S \in \mathcal{S}'} S \right| = \sum_i \sum_{S \in \mathcal{S}'_i} |S| \leq \sum_i (4s^2i)^{2s+1} \cdot c^{s-\ell} \cdot 2^{-i}|C_{\ell-1}|.$$

By Lemma 4.1, we have

$$\sum_i (4s^2i)^{2s+1} 2^{-i} = (2s)^{4s+2} \sum_i i^{2s+1}/2^i \leq (2s)^{4s+2} \cdot 2(2s+1)^{2s+1} < s^{15s}/5 = c/5.$$

Therefore, $\left| \bigcup_{S \in \mathcal{S}'} S \right| < \frac{c^{s-\ell+1}}{5} \cdot |C_{\ell-1}|$, which is a contradiction. Thus there exists an index i for which $|\mathcal{S}'_i| > (4s^2i)^{2s+1}$.

Let $C = (a_1, \dots, a_{d(C_{\ell-1})+i}, *, \dots, *)$ and consider it as a level ℓ cube. By Proposition 4.2, there are at most $s \cdot i$ cubes in \mathcal{C} with codimension at most $d(C)$ which have level ℓ adjacency with C (remember that these may have level higher than ℓ). Let $\mathcal{A}^{(\ell)}$ be the family consisting of these cubes and note that the dominating parameter of C and A for $A \in \mathcal{A}^{(\ell)}$ is always i . For a cube $A \in \mathcal{A}^{(\ell)}$, we say that a set $S \in \mathcal{S}'_i$ is *bad* for A if there are at least $(4s^2i)^{\ell(S_A)+\ell(S)-2s}|S_A||S| = (4s^2i)^{\ell(S_A)+\ell-2s}|S_A||S|$

blue edges between S_A and S . Otherwise, we say that S is *good* for A . We claim that there are at most $(4s^2i)^{2s}$ sets in \mathcal{S}'_i which are bad for each fixed A .

Let $X_i = \bigcup_{S \in \mathcal{S}'_i} S$ and note, by Proposition 4.1, that there are at most $|S_A| \cdot 2c^{s-\ell} \cdot 2^{n-d(C)}$ blue edges between the sets S_A and X_i . Each set $S \in \mathcal{S}'_i$ which is bad for A accounts for at least $(4s^2i)^{\ell(S_A)+\ell-2s} |S_A| |S| \geq \frac{c^{s-\ell} \cdot 2^{n-d(C)}}{(4s^2i)^{2s-1}} |S_A|$ such blue edges (note that $|S| = c^{s-\ell} \cdot 2^{n-d(C)}$ and $\ell \geq 1$). Therefore, in total, there are at most $(4s^2i)^{2s}$ sets in \mathcal{S}'_i which are bad for A , as claimed above. Since there are at most si cubes in $\mathcal{A}^{(\ell)}$ and $|\mathcal{S}'_i| > (4s^2i)^{2s+1}$, there exists a set $S \in \mathcal{S}'_i$ which is good for all the cubes $A \in \mathcal{A}^{(\ell)}$.

In order to show that C satisfies (a) and (b), it suffices to verify that $d(C) \geq d$, since this implies the fact that C is disjoint from all the other cubes of level at least ℓ (note that $C \subset C_{\ell-1}$ and if C intersects some other cube of level at least ℓ , then that cube must contain C and therefore also contains (a_1, \dots, a_n) by Proposition 2.2). Furthermore, if this is the case, assigning C to S is a proper assignment (thus we have (c)).

Now suppose, for the sake of contradiction, that $d(C) < d$ and consider the time t immediately after we last embedded a cube of codimension at most $d(C)$. At time t , since $d(C) \geq d(C_{\ell-1})$, the cube $C_{\ell-1}$ was already embedded and, since there are no cubes of level ℓ covering (a_1, \dots, a_n) , the cubes in \mathcal{C}_ℓ are disjoint from C . For $\rho \leq \ell-1$, a cube in the partial tiling at time t , which is of codimension at most $d(C)$, is ρ -adjacent to C if and only if it is ρ -adjacent to C_ℓ . Hence, the family of cubes which are adjacent to C at time t is a subfamily of $\bigcup_{\rho=1}^\ell \mathcal{A}^{(\rho)}$. Therefore, C could have been added to the tiling at time t as well, contradicting the fact that we always choose a cube of minimum codimension. Thus we have $d(C) \geq d$, as claimed. \square

Note that as an outcome of our algorithm, we obtain a tiling \mathcal{C} such that for every pair of adjacent cubes $C, C' \in \mathcal{C}$, we have control on the number of blue edges between S_C and $S_{C'}$ (as given in the definition of proper assignment).

4.3 Imposing a maximum degree condition

As in the previous sections, we now impose certain maximum degree conditions between the sets S_C for $C \in \mathcal{C}_{s-2}$. For a set $C \in \mathcal{C}_{s-2}$ of codimension $d = d(C)$ and relative codimensions $d_\ell = d_\ell(C)$, $1 \leq \ell \leq s-2$, recall that we have a set $S_C \in \mathcal{S}$ such that $|S_C| \geq c^2 \cdot 2^{n-d}$. Let $\mathcal{A}^{(\rho)}$ be the family of cubes in \mathcal{C}_{s-2} with codimension at most d which have level ρ adjacency with C . By Proposition 4.2, we have $|\mathcal{A}^{(\rho)}| \leq d_\rho(C)$ for each $\rho = 1, \dots, s-2$. For each $A \in \mathcal{A}^{(\rho)}$, let $\delta_A = \delta(C, A)$ and note that $\delta_A = \max\{d_\rho(A), d_\rho(C)\} \geq d_\rho(C)$. Also, since $\ell(A) + \ell(C) - 2s = -4$, there are at most $\frac{1}{(4s^2\delta_A)^4} |S_C| |S_A|$ blue edges between S_C and S_A .

Now for $\rho = 1, \dots, s-2$, and each $A \in \mathcal{A}^{(\rho)}$, remove all the vertices in S_C which have at least $\frac{1}{4s^2\delta_A} |S_A|$ blue neighbors in S_A and let T_C be the subset of S_C left after these removals. Since there are at most $\frac{1}{(4s^2\delta_A)^4} |S_C| |S_A|$ blue edges between S_C and S_A , we remove at most $\frac{|S_C|}{(4s^2\delta_A)^3}$ vertices from S_C for each set $A \in \mathcal{A}^{(\rho)}$. Thus the resulting set T_C is of size at least

$$|T_C| \geq |S_C| - \sum_{\rho=1}^{s-2} d_\rho \cdot \frac{|S_C|}{(4s^2\delta_A)^3} \geq \frac{|S_C|}{2} \geq c \cdot 2^{n-d}.$$

For each $A \in \mathcal{A}^{(\rho)}$, all the vertices in T_C have blue degree at most $\frac{1}{4s^2\delta_A} |S_A| \leq \frac{1}{2s^2\delta_A} |T_A|$ in the set

T_A . Thus we obtain the following property.

Maximum degree condition. Let $C, C' \in \mathcal{C}_{s-2}$ be a pair of cubes having level ρ adjacency with $d(C) \geq d(C')$. Then every vertex in T_C has at most $\frac{1}{2s^2\delta(C, C')}|T_{C'}|$ blue neighbors in the set $T_{C'}$.

4.4 Embedding the cube

We now show how to embed Q_n . Recall that we found an $(s-1)$ -tiling \mathcal{C} of Q_n . We will greedily embed the cubes in the level $s-2$ tiling \mathcal{C}_{s-2} one by one into their assigned sets from the family \mathcal{S}_{s-2} , in decreasing order of their codimensions. If there are several cubes of the same codimension, then we arbitrary choose the order between them.

Suppose that we are about to embed the cube $C \in \mathcal{C}_{s-2}$. Let $d = d(C)$ and $d_\ell = d_\ell(C)$, for $\ell = 1, \dots, s-2$. We will greedily embed the vertices of C into $T_C \subseteq S_C$. Suppose that we are about to embed $x \in C$ and let $f : Q_n \rightarrow [N]$ denote the partial embedding of the cube Q_n obtained so far. For each ρ , let A_ρ be the set of neighbors of x which are already embedded and belong to a cube other than C that has level ρ adjacency with C . Note that we have $|A_\rho| \leq d_\rho$ for every ρ . Since we have so far only embedded cubes of codimension at least d , for each ρ , the vertices $f(v)$ for $v \in A_\rho$ have blue degree at most $\frac{1}{2s^2d_\rho}|T_C|$ in the set T_C , by the maximum degree condition imposed in Section 4.3. Together, these neighbors forbid at most

$$\sum_{\rho=1}^{s-2} d_\rho \cdot \frac{1}{2s^2d_\rho}|T_C| \leq \frac{1}{2}|T_C|.$$

vertices of T_C from being the image of x .

In addition, x has at most $n-d$ neighbors which are already embedded and belong to C . By Proposition 4.1, for each such vertex v , $f(v)$ has blue degree at most $\frac{2^{n-d}}{n}$ in the set T_C . Together, these neighbors forbid at most 2^{n-d} vertices of T_C from being the image of x . Finally, there are at most $2^{n-d} - 1$ vertices in T_C which are images of some other vertex of C that is already embedded. Therefore, the number of vertices in T_C into which we cannot embed x is at most

$$\frac{1}{2}|T_C| + 2^{n-d} + (2^{n-d} - 1) < |T_C|,$$

where the inequality follows since $|T_C| \geq c \cdot 2^{n-d}$. Hence, there exists a vertex in T_C which we can choose as an image of x to extend the current partial embedding of the cube. Repeating this procedure until we embed the whole cube Q_n completes the proof.

5 Small separators, forbidden minors and Ramsey goodness

Let \mathcal{G}_H be the family of graphs G which do not contain an H -minor. In order to show that this family is s -good, we wish to apply the following result of Nikiforov and Rousseau [30].

Theorem 5.1. *For every $s \geq 3$, $d \geq 1$ and $0 < \gamma < 1$, there exists $\eta > 0$ such that the class \mathcal{G} of d -degenerate graphs G with a $(|V(G)|^{1-\gamma}, \eta)$ -separator is s -good.*

There are two conditions here that need to be verified in order to gain the conclusion of Theorem 1.2. Firstly, we need to show that the graphs in \mathcal{G}_H have bounded degeneracy. This was first proved by Mader [28]. Later, Kostochka [24, 25] and Thomason [35] independently established the following bound, which is tight apart from the constant factor. More recently, the asymptotic value of c was determined by Thomason [36].

Theorem 5.2. *There exists a constant $c > 0$ such that any graph with average degree at least $ch\sqrt{\log h}$ contains a K_h -minor.*

Secondly, we need to show that the graphs in \mathcal{G}_H have appropriate separators. We will use the following result of Alon, Seymour and Thomas [2]. Note that, in the particular case of planar graphs, such a separator theorem was proved much earlier by Lipton and Tarjan [27].

Theorem 5.3. *Let G be a graph on n vertices containing no K_h -minor. Then G has an $(h^{3/2}n^{1/2}, \frac{2}{3})$ -separator.*

We are now ready to prove our main result about \mathcal{G}_H , which we recall from the introduction.

Theorem 5.4. *For every fixed graph H , the class \mathcal{G}_H of graphs G which do not contain an H -minor is s -good for all $s \geq 3$.*

Proof. In order to apply Theorem 5.1, it is enough to verify that, for any $\eta > 0$, any sufficiently large graph G in \mathcal{G}_H has bounded degeneracy and a $(|V(G)|^{2/3}, \eta)$ -separator.

Suppose that the graph H has h vertices. Note, by Theorem 5.2, that any graph with average degree at least $ch\sqrt{\log h}$ contains a K_h -minor and, hence, an H -minor. This implies that the average degree of every subgraph of G is at most $ch\sqrt{\log h}$. In turn, this easily implies that the degeneracy of G is at most $ch\sqrt{\log h}$.

Suppose that G has n vertices. To show that a sufficiently large graph G from the class \mathcal{G}_H contains an $(n^{2/3}, \eta)$ -separator T , we begin by applying Theorem 5.3 to conclude that there is an $(h^{3/2}n^{1/2}, \frac{2}{3})$ -separator. Note, by taking unions of small components if necessary, that this gives a decomposition of the vertex set of G into three sets T, A and B such that $|T| \leq h^{3/2}n^{1/2}$, $|A|, |B| \leq \frac{2}{3}n$ and there are no edges between A and B .

We will prove, by induction on i , that for n sufficiently large depending on i , there is a separator T_i of size at most $2^i h^{3/2} n^{1/2}$ that splits the vertex set of G into 2^i sets $U_{i,1}, \dots, U_{i,2^i}$, each of size at most $(\frac{2}{3})^i n$, so that there are no edges between any distinct sets $U_{i,a}$ and $U_{i,b}$. By the previous paragraph, the result holds for $i = 1$. Now suppose that it holds for i . We will show that a similar conclusion follows for $i + 1$.

To begin, we apply Theorem 5.3 within each of the sets $U_{i,j}$ to conclude that there is a decomposition of $U_{i,j}$ into sets $T_{i,j}, A_{i,j}$ and $B_{i,j}$ such that $|T_{i,j}| \leq h^{3/2}|U_{i,j}|^{1/2}$,

$$|A_{i,j}|, |B_{i,j}| \leq \frac{2}{3}|U_{i,j}| \leq \left(\frac{2}{3}\right)^{i+1} n$$

and there are no edges between $A_{i,j}$ and $B_{i,j}$. We let the collection $\{U_{i+1,j}\}_{j=1}^{2^{i+1}}$ consist of all sets of the form $A_{i,j}$ and $B_{i,j}$ for $j = 1, \dots, 2^i$. This collection has size 2^{i+1} , each of the sets has size at

most $\left(\frac{2}{3}\right)^{i+1} n$ and there are no edges between distinct sets $U_{i+1,a}$ and $U_{i+1,b}$. If we also let T_{i+1} be the union of T_i and $T_{i,j}$, for $1 \leq j \leq 2^i$, we have

$$|T_{i+1}| \leq |T_i| + \sum_{j=1}^{2^i} |T_{i,j}| \leq 2^i h^{3/2} n^{1/2} + \sum_{j=1}^{2^i} h^{3/2} |U_{i,j}|^{1/2} \leq 2^i h^{3/2} n^{1/2} + 2^i h^{3/2} n^{1/2} = 2^{i+1} h^{3/2} n^{1/2}.$$

Therefore, the induction holds.

If we now apply this result with $i = 2 \log \eta^{-1}$ and $n \geq h^9 \eta^{-12}$, we see, since $\left(\frac{2}{3}\right)^i \leq \eta$ and $n^{2/3} \geq 2^i h^{3/2} n^{1/2}$, that G has an $(n^{2/3}, \eta)$ -separator, as required. The result follows. \square

Let \mathcal{K} be the collection of graphs K for which there is a proper vertex coloring in $\chi(K)$ colors such that at least two of the color classes have size one. The full result of Nikiforov and Rousseau (namely, Theorem 2.2 of [30]) says that for any $K \in \mathcal{K}$, $d \geq 1$ and $0 < \gamma < 1$ there exists $\eta > 0$ such that the class \mathcal{G} of d -degenerate graphs G with a $(|V(G)|^{1-\gamma}, \eta)$ -separator is K -good. This may in turn be used to show that for any H the family of graphs \mathcal{G}_H is K -good for all $K \in \mathcal{K}$. However, \mathcal{G}_H is not K -good for all graphs K . This follows from the observation mentioned in the introduction that $K_{1,t}$ is not $K_{2,2}$ -good for any t .

The family of graphs \mathcal{G}_H with forbidden H -minor is not the only class of graphs known to have small separators. For example, several geometric separator theorems are known (see, for example, [19, 29]) saying that the class of intersection graphs formed by certain collections of bodies have small separators. In any of these cases, Theorem 5.1 will also allow us to show that the classes are K -good for any $K \in \mathcal{K}$, though in some cases we may have to further restrict the class in order to obtain the required degeneracy condition.

6 Concluding remarks

The question of determining whether the cube is s -good for any $s \geq 3$ is only one of two well-known questions of Burr and Erdős regarding Ramsey numbers and the cube. The other [10, 14] is the question of determining whether the Ramsey number $r(Q_n) := r(Q_n, Q_n)$ of the cube with itself is linear in the number of vertices in Q_n .

Beginning with Beck [3], who proved that $r(Q_n) \leq 2^{cn^2}$, a large number of papers [22, 23, 26, 32, 33] have considered this question, with the current best bound being $r(Q_n) \leq n^{2^{2n+5}}$, due to Fox and Sudakov [20]. That is, $r(Q_n) \leq |Q_n|^{2+o(1)}$. This (almost) quadratic bound for $r(Q_n)$ was obtained using a careful application of dependent random choice [21]. It seems likely that in order to improve the bound one will have to improve this latter technique. However, we do not rule out the possibility that some of the embedding lemmas used in the current paper could also be of use.

Suppose H is an N -vertex graph with chromatic number r , maximum degree $O(1)$ and bandwidth $o(N)$. Bollobás and Komlós conjectured that every N -vertex graph G which does not contain H as a subgraph has minimum degree at most $(1 - \frac{1}{r} + o(1))N$. This conjecture was recently verified by Böttcher, Schacht and Taraz [5]. It is natural to wonder to what extent the bounded maximum degree condition on H can be relaxed. In particular, is it true that every Q_n -free graph on 2^n vertices has minimum degree at most $(\frac{1}{2} + o(1))2^n$? Note that Q_n is a bipartite graph on $N = 2^n$ vertices, which is n -regular with $n = \log N$, and has bandwidth $O(N/\sqrt{\log N})$. It appears that new techniques would have to be developed to handle this problem. The proof of the Bollobás-Komlós

conjecture uses the regularity lemma and the blow-up lemma, giving quantitative estimates that are too weak to embed spanning subgraphs which are as dense as cubes. A positive solution would likely lead to a proof of the Burr-Erdős conjecture discussed above that cubes have linear Ramsey number.

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Note added in proof. Recently, the main conjecture studied in this paper, that $r(Q_n, K_s) = (s-1)(2^n-1)+1$ for s fixed and n sufficiently large, was resolved by Fiz Pontiveros, Griffiths, Morris, Saxton and Skokan [17, 18]. Their proof builds upon the techniques developed in this paper.

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