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On Some Edge Rotation Distance Graphs

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Abstract: The concept of edge rotations and distance between graphs was introduced by Gary Chartrand et al [1]. A graph G can be transformed into a graph H by an edge rotation if G contains distinct vertices u, v and w such $uv \in E(G)$ and $uw \notin E(G)$ and $H \cong G - uv + uw$. In this case, G is transformed into H by "rotating" the edge uv of G into uw . In this paper we consider rotations on generalized Petersen graphs and minimum self-centered graphs. We have also developed algorithms to generate distance degree injective (DDI) graphs and almost distance degree injective (ADDI) graphs from cycles using the concept of rotations followed by some general results.

Keywords: ADDI graphs, average eccentricity, DDI graphs, edge rotations, edge rotation distance graphs, Generalized Petersen graph, r -distance graph.

I. Introduction

Unless mentioned otherwise, for terminology and notation the reader may refer to Buckley and Harary [23] and Chartrand and Zang [24], new ones will be introduced as and when found necessary.

In this paper, by a graph G , we mean a simple, undirected, connected graph without self-loops. The order and size are respectively the number of vertices denoted by n and number of edges denoted by m .

The distance $d(u, v)$ between any two vertices u and v , of G , is the length of the shortest path between u and v . The eccentricity $e(u)$ of a vertex u is the distance to a farthest vertex from u . The maximum and the minimum eccentricity amongst the vertices of G are respectively called the diameter $diam(G)$ and radius $rad(G)$. If $diam(G) = rad(G)$, then the graph G is said to be self-centered graph. If $d(u, v) = e(u)$, ($v \neq u$) then we say that v is an eccentric vertex of u .

The distance degree sequence (dds) of a vertex in a graph $G = (V, E)$ is the list of number of vertices at distance $1, 2, \dots, e(v)$ in that order, where $e(v)$ denotes the eccentricity of v . Thus the sequence $(d_{i_0}, d_{i_1}, d_{i_2}, \dots, d_{i_j})$, is the distance degree sequence (dds) of the vertex v_i in G where d_{i_j} denotes the number of vertices at distance j from v_i .

The concept of distance degree regular (DDR) graphs was introduced by Bloom et al. [26] as the graph for which all vertices have the same distance degree sequence. This was further studied by Bloom et al. [25], Halberstam et al. [27], Itagi Huilgol et al. [2], [3], [4].

The other extreme of distance degree regular (DDR) graphs is the distance degree injective (DDI) graphs. The concept of distance degree injective graphs was introduced by Bloom et al. in [25]. A graph G is said to be a distance degree injective (DDI) graph if no two of its vertices have the same distance degree sequence. In literature, characterizations of both DDR and DDI graphs are not known. But there are several particular cases, [26], [27], [2], [3], [4], [28], etc. So the construction of distance degree injective graphs or the almost distance degree injective graphs (DDI/ADDI) is also a challenging one. In [3] Itagi Huilgol et al. have constructed higher order DDI graphs using products. Fast generation of cubic graphs was done by Brinkmann [5]. In [6] Itagi Huilgol et al. have introduced the concept of Almost Distance degree Injective (ADDI) graphs. A graph G of order n is said to be ADDI or almost DDI if $n - 2$ vertices have different distance degree sequence and two vertices have the same distance degree sequence.

The concept of distance between isomorphism classes of graphs was introduced by Zelinka [7] which was later extended to trees [8] also. Based on these two papers the concept of rotations called 'edge rotations' and the distance between such graphs was introduced by Chartrand et al. [1]. A graph G can be transformed into a graph H by an edge rotation given by $H \cong G - uv + uw$ where u, v and w are distinct vertices of G such that $uv \in E(G)$ and $uw \notin E(G)$. Later in [9] Zelinka gave a comparison of various distance for the isomorphism classes of graphs and trees, which was based on the concept of edge rotations. It was also showed that this distance is a metric by Balaz et al. in [10].

Zelinka studied various aspects using the concept of distance between graphs and edge rotations in [11], [12] and [13]. The concept of "edge move" was defined by Balaz et al. [10] and is defined in this manner. A graph G can be transformed into a graph H by an edge move given by $H \cong G - uv + wx$

$uv \in E(G)$ and $xw \in E(\overline{G})$. Later Johnson in [14] introduced a new kind of edge rotation called the “edge shift”. It is defined as an edge rotation for the given graph A , $t = (u, v, w)$ such that vw is an edge of A . As with an edge rotation, tA will denote the newly formed graph $A - uv + uw$. This new edge rotation is again a metric.

Later an inequality between edge shift, edge rotation and edge move was proposed by Benade et al. [15]. The upper and the lower bounds were given by Faudree et al. in [16]. Under the concept of labeling, the distances between graphs using edge operation was done by Goddard et al. in [17].

The rotation distance between graphs G and H is denoted by $d_r(G, H)$, if there exists a sequence of graphs G_1, G_2, \dots, G_{k-1} such that G_1 is obtained by an edge rotation on G , and for each $1 \leq i \leq k-1$, G_{i+1} is obtained by an edge rotation on G_i , with H obtained from G_{k-1} by one edge rotation. In this case we denote the rotation distance from G to H as $d_r(G, H)$ and it is equal to k .

Definition 1: [1] Let $S = G_1, G_2, \dots, G_k$ be a set of graphs all of same order and the same size. Then the rotation distance graph $D(S)$ of S has S as its vertex set and vertices (graphs) G_i and G_j are adjacent if $d_r(G_i, G_j) = 1$, where $d_r(G_i, G_j)$ is the rotation distance between G_i and G_j .

A graph G is a *edge rotation distance graph (ERDG) (or r-distance graph)* if $G \cong D(S)$ for some set of graphs.

A study done by Chartrand et al. [18] showed that the cycles, the complete bipartite graph $K_{3,3}$ and $K_{2,p}$ ($p \geq 1$) are *edge rotation distance graphs (ERDG)*. This was later extended by Jarrett [19] in showing that complete graphs, trees and the wheel $(W_{1,n})$ are also edge rotation distance graphs. Jarrett also gave a different proof other than the one due to Chartrand et al. [18]. It was shown that the complete bipartite graph $K_{m,n}$ ($3 \leq m \leq n$) is a edge rotation distance graph.

In this paper we consider the edge rotations for the generalized Petersen graphs $(G_p(n, k))$ and show that it is an edge rotation distance graph for $n \geq 5$, $n \in \mathbb{Z}^+$, $k = 1$ where \mathbb{Z}^+ is the set of all positive integers. We have also developed two algorithms which show the construction of distance degree injective graphs and almost distance degree injective graphs from cycles using the concept of edge rotation, some results on average eccentricity with respect to edge rotation are also proved.

II. Edge Rotations for Generalized Petersen Graphs $G_p(n, k)$

The generalized Petersen graphs were introduced by Coxeter in [20] and later named by Watkins in [21].

Definition 2: For integers n and k with $2 \leq 2k \leq n$, the generalized Petersen graph $G(n, k)$ has the vertex set $V(G(n, k)) = u_0, u_1, \dots, u_{n-1}, v_0, v_1, \dots, v_{n-1}$ and the edge set $E(G(n, k)) = [u_i, u_{i+j}], [u_i, v_i], [v_i, v_{i+k}]$, where i is an integer and all subscripts are read modulo n .

As the name suggests $G(n, k)$ is generalization for the Petersen graph. In particular for $n = 5$, $k = 2$, we get the Petersen graph. Note that a generalized Petersen graph is a cubic graph.

In this section we consider edge rotations on the generalized Petersen graph where $n \geq 5$, $n \in \mathbb{Z}^+$, $k = 1$.

The new graph obtained after a single edge rotation will be called as G' , i.e., $G' = G - e(G) + e(\overline{G})$

Theorem 2.1: An edge rotation in $G_p(n, 1)$ when n is odd results in a self-centered graph if the rotation of any edge induces a cycle of length 5 or 3.

Proof: By the definition of rotation: $G' = G - e(G) + e(\overline{G})$ where G' is the new graph obtained after a single edge rotation. As $G_p(n, 1)$ is cubic graph, on rotation, a change in degrees of two vertices resulting in a non-regular graph. The girth of any generalized Petersen graph is always four. Hence upon rotation, G' induces a cycle of minimum length three to a maximum length $(n-1)$. Among these, similar to the rotations shown in “Fig. 1”, result in a self-centered graph.

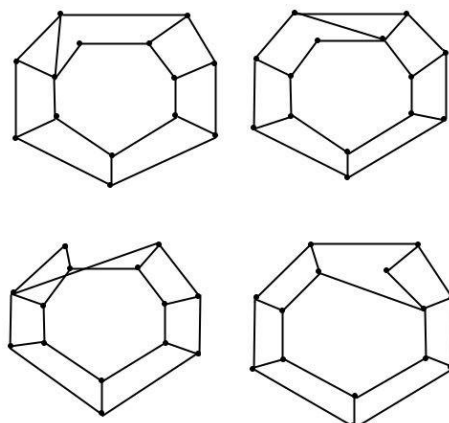


Fig 1: Rotations on a $G_p(7,1)$ graph

Remark 1: An edge rotation made on a $G_p(n,1)$ when n is even does not allow the graph to be self-centered.

Remark 2: An edge rotation on $G_p(3, 1)$ retains the radius thus resulting in a self-centered graph.

Remark 3: The complement of a circulant graph is not a Edge Rotation Distance Graph if the circulant consists of more than two jump sizes.

It was shown by E. B. Jarrett [19] that the cycles (for $n \geq 3$), the complete bipartite graph $K_{m, n}$ ($3 \leq m \leq n$) and the wheel $(W_{1, n})$ is a ERDG. Now by a slight modification of the graphs used in proofs of the theorems proved by Jarrett in [19], we show that $G_p(n, k)$ where $k = 1$ result in ERDG's. We first generate two cycles as in [19] and then show that the rotation distance graph on this class of cycles results in a generalized Petersen graph.

Theorem 2.2: The generalized Petersen graph $G_p(n, 1)$ where $n \geq 3$ and $n \in N$ is a edge rotation distance graph(ERDG).

Proof: Let P be a path of length $3n - 1$ where n is a natural number. Let the vertices of the path be denoted by $P: v_0, v_1, v_2, \dots, v_{2n+1}, \dots, v_{3n-1}$. Let G be a graph obtained from P by adding two new vertices a_1, a_2 and three new edges $v_{2n+1}a_1, a_1a_2, a_2v_{2n+1}$. Then for $i = 1, 2, \dots, n - 1$, we define G_i to be the graph obtained from G by adding two new vertices x and y such that x is adjacent to v_{2i} and y is adjacent to v_{2i+2} . We also define G_n as the graph obtained from G by adding two new vertices namely x adjacent only to v_{2n} and vertex y adjacent only to v_2 .

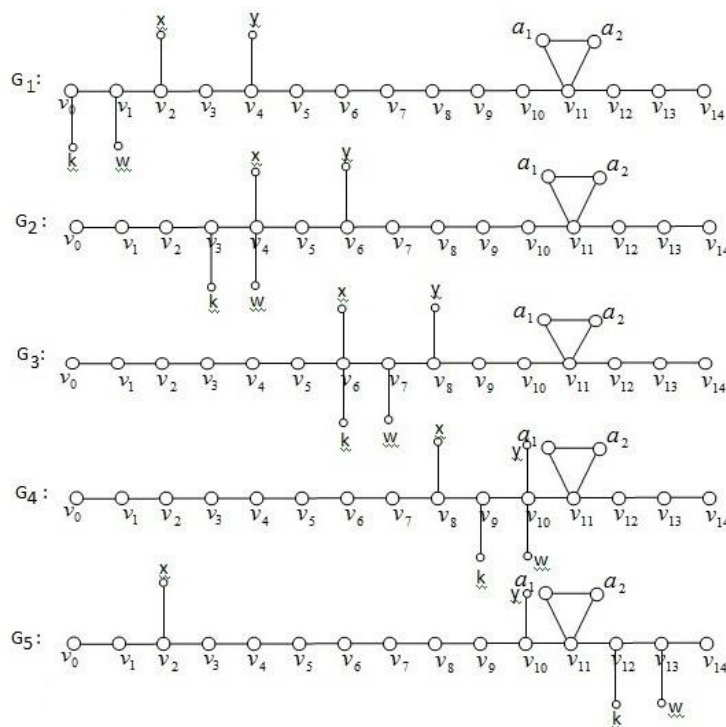


Fig 2: Rotation on graph G

Since an edge rotation changes the degrees of exactly two vertices $d_r(G_i, G_j) > 1$ for integers i and j , such that $1 \leq i \leq j \leq n$. On the other hand, for $i = 1, 2, \dots, n - 2$, the graph $G_{i+1} \cong G_i - xv_{2i} + xv_{2i+4}$ and consequently $d_r(G_i, G_{i+1}) = 1$. Similarly $d_r(G_{n-1}, G_n) = d_r(G_n, G_1) = 1$, since $G_n \cong G_{n-1} - xv_{2n-2} + xv_2$ and $G_1 \cong G_{n-1} - xv_{2n} + xv_4$; thus $d_r(G_1, G_2, G_3, \dots, G_n) \cong C_n$.

In a similar way we shall construct one more graph H and consider the path to consists of vertices $P_1: v_0, v_1, v_2, \dots, v_{2n+1}, \dots, v_{3n-1}$. Let G be a graph obtained from P_1 by adding two new vertices a_1, a_2 and three new edges $v_{2n+1}a_1, a_1a_2, a_2v_{2n+1}$. Then for $i = 1, 2, \dots, n - 1$, we define H_i to be the graph obtained from H by adding two new vertices x and y such that x is adjacent to v_{2i} and y is adjacent to v_{2i+2} . We also define H_n as the graph obtained from H by adding two new vertices namely x adjacent only to v_{2n} and vertex y adjacent only to v_2 .

We now show that the edge rotation distance between each of these G_i and H_i is equal to one. To show that this relation exists we add two new vertices namely w and k and then rotate these edges such that $d_r(G_i, H_i) = 1$, for $i = 1, 2, 3, \dots, n$.

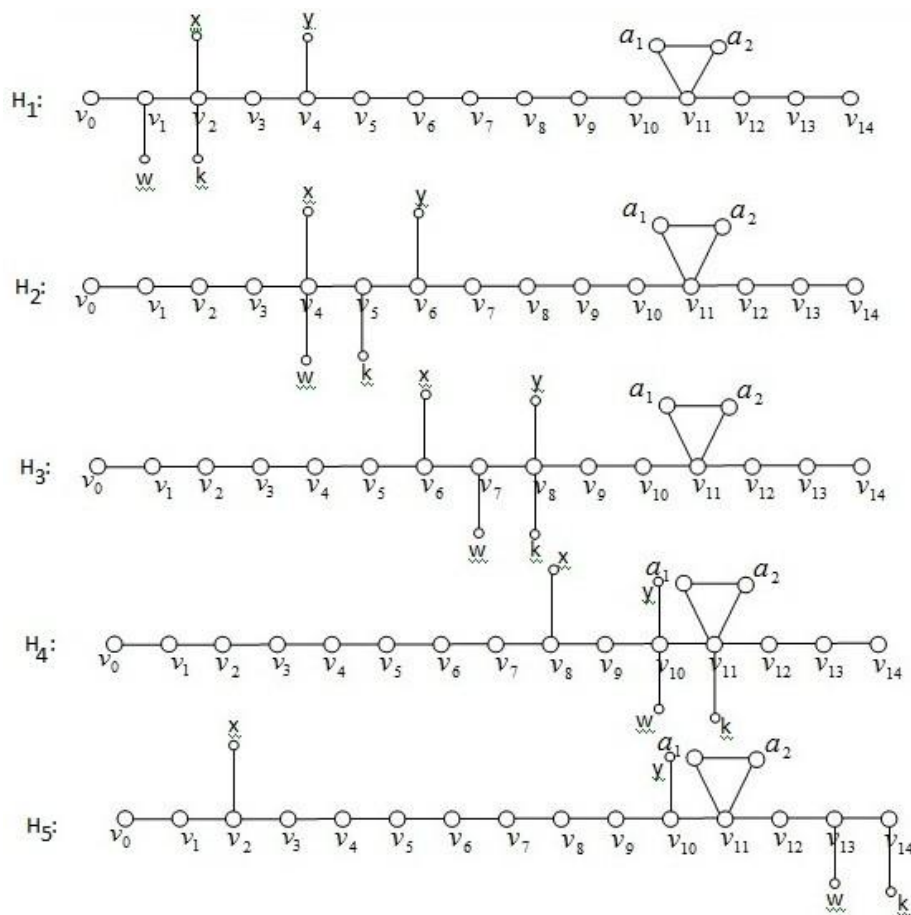


Fig 3: Rotations on graph H

For $i = 1, 2, \dots, n$, we first define G_i to be the graph obtained from G by adding two other vertices 'k' adjacent to v_0 and 'w' adjacent to v_1 . The next graph G_2 is obtained by adding the vertices k and w adjacent to v_3 and v_4 . In G_3 , k is adjacent to v_6 and w is adjacent to v_7 . The other G_i 's are obtained in a similar pattern up to $i = n - 1$. In the last graph when $i = n$, the vertices k and w are adjacent to $3n - 3$ and $3n - 2$.

For $i = 1, 2, \dots, n$, we first define H_1 to be the graph obtained from H by adding two other vertices 'w' adjacent to v_1 and 'k' adjacent to v_2 . The next graph H_2 is obtained by adding the vertices w and k adjacent to v_4 and v_5 . H_3 is obtained by making w adjacent to v_7 and k adjacent to v_8 . The other H_i 's are obtained in a similar pattern up to $i = n - 1$. In the last graph when $i = n$, the vertices w and k are adjacent to $3n - 3$ and $3n - 2$.

Since an edge rotation changes degrees of exactly two vertices $d_r(G_b, G_i) > 1$. On the other hand for $i = 1, 2, 3, \dots, n$, $H_1 \cong G_1 - kv_0 + kv_2$, $H_2 \cong G_2 - kv_3 + kv_5$ and so on. Consequently, $d_r(G_b, H_i) = 1$. Hence, the rotation distance between each of G_i and H_i is always one.

We consider an example to generate a generalized Petersen graph for $n = 5$ and give an algorithmic construction of the graph which can be applied to higher order generalized Petersen graph.

Example 2.3 Consider the graph $G_p(5, 1)$

STEP - 1: For the graphs G_i and H_i obtained after adding the vertices a_1, a_2 which are used in the generation of the cycle, we add two new vertices namely 'k' and 'w' to G_i and H_i . The adjacencies between these vertices is establishes as follows.

- Join 'k' to v_0 and 'w' to v_1 in G_1 . Now in H_1 , join 'w' to v_1 and 'k' to v_2 .
- In G_2 , join 'k' to v_3 and 'w' to v_4 . In H_2 , join 'w' to v_4 and 'k' to v_5 .
- In G_3 , join 'k' to v_6 and 'w' to v_7 . In H_3 , join 'w' to v_7 and 'k' to v_8 .
- In G_4 , join 'k' to v_9 and 'w' to v_{10} . In H_4 , join 'w' to v_{10} and 'k' to v_{11} .
- In G_5 , join 'k' to v_{12} and 'w' to v_{13} . In H_5 , join 'w' to v_{13} and 'k' to v_{14} .

STEP - 2: Apply the concept of edge rotation to the newly added vertices.

STEP – 3: We now rotate the newly added edges between each of G_i and H_i . That is the edge v_0k in G_1 is rotated to v_2k in H_1 and thus $d_r(G_1, H_1) = 1$.

In a similar way the rotation is carried in the rest of the graphs to establish the r-distance relation.

Also, $d_r(G_2, H_2) = d_r(G_3, H_3) = d_r(G_4, H_4) = d_r(G_5, H_5) = 1 = d_r(G_1, H_1)$.

STEP – 4: Thus, a relation is brought between the vertices of G and H in showing that the generalized Petersen graph is a ERDG.

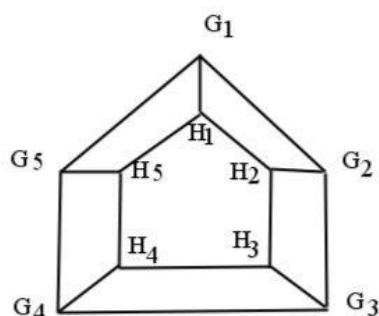


Fig 4: Edge rotation on $G_p(5,2)$

III. Edge Rotation on $K_{1,m}$

Here we prove that the star $K_{1,m}$ is a edge rotation distance graph. The proof is an algorithmic one.

Lemma 3.1:The star $K_{1,m}$, where $m \in \mathbb{Z}^+$ is a ERDG (for $m \geq 2$).

Proof: Consider the path P of order $l + m$, where $l + m$ is the order of the star considered. We give a step by step procedure as proof to show that the star is a ERDG.

STEP – 1: Let $P = v_1, v_2, \dots, v_{m+1}$ be a path and G be the graph formed by adding two new vertices v_{m+2}, v_{m+3} .

STEP – 2: Join the vertices $v_{m+1} v_{m+2}, v_{m+2} v_{m+3}, v_{m+3} v_{m+1}$.

STEP – 3: Add a new vertex “ x ” to $G_i, i = 1, 2, \dots, m+1$ such that x is adjacent to v_i in G_i .

STEP – 4: $d_r(G_1, G_{1+k}) = 1, k = 1, 2, \dots, m$.

That is $d_r(G_1, G_2) = d_r(G_1, G_3) = \dots = d_r(G_1, G_m) = 1$.

Here we observe that $\deg(v_1)$ in $G_1 = 2$. Similarly $\deg(v_2)$ in $G_2 = 3, \deg(v_3)$ in $G_3 = 3$. The presence of a degree two vertex in G_1 at v_1 changes the degree of that vertex upon rotation. There exists at least one vertex of degree three in each of G_2 to G_{m+1} which is the result of edge rotation.

STEP – 5 : Thus the star $K_{1,m}$ is a ERDG.

Example 3.1:The following example shows that $K_{1,4}$ is a ERDG. Similar construction can be given to any positive integer m .

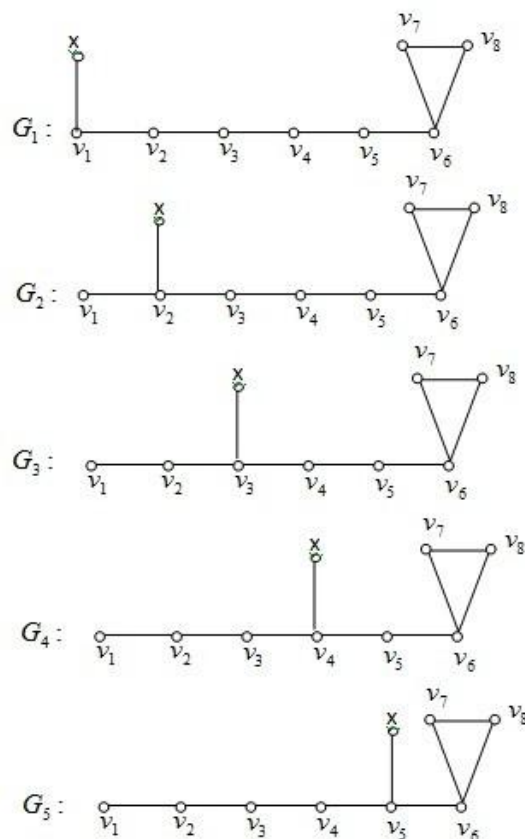


Fig 5: The star $K_{1,m}$ is a ERDG.

Remark 4: Let G be a graph with a single cut – edge, then the graph becomes disconnected if the cut- edge is rotated to any other vertex on the same component.

Definition 3: [23] The eccentric mean or average eccentricity is defined as $\mu_e(G) = \left(\frac{1}{n}\right) \sum_{i=1}^n m_i e_i$ taken over all vertices in the graph where m_i 's are the multiplicities of the eccentricities e_i .

Theorem 3.2 : For any even cycle C_n , for $n > 4$, then average eccentricity of C'_n lies between $\left[\frac{n}{2} - 1, \frac{n}{2} + 1\right]$ for any edge rotation in the cycle.

Proof: Let $u_1, u_2, \dots, u_n, u_1$ be an even cycle. Let u_1u_2 be the edge to be rotated to any other vertex say u_3, u_4, \dots, u_{n-1} or u_2u_1 be the edge to be rotated to any other vertex say u_4, u_5, \dots, u_n . The rotation made induces a cycle of minimum length three to a maximum length $n - 1$. As their length of the cycle induced increases, the eccentricity of the vertices also increases. We observe that there exist at least two or more vertices which have the same eccentricity after the edge rotation. This occurs at the vertex where the rotation has been performed since this is the only vertex where we find two or more vertices with the same eccentricity and the rest of the vertices with different eccentricities. Thus we find that the average eccentricity of the graph C'_n reduces by one upon rotation. Hence the eccentricity ranges from a minimum of $\frac{n}{2} - 1$ to a maximum of $\frac{n}{2} + 1$ upon different rotations. Thus when the average is taken over all the vertices we find that the average eccentricity lies between the closed interval $\left[\frac{n}{2} - 1, \frac{n}{2} + 1\right]$.

Lemma 3.2 For any odd cycle C_n the average eccentricity of C'_n is not greater than $\left\lfloor \left(\frac{n}{2}\right) \right\rfloor + 2$.

Proof : A single edge rotation induces a cycle of minimum length three and a maximum length $n-1$. The graph C'_n formed is a cycle followed by a path of some length “ m ”. The eccentricity of each vertex ranges from $(n/2)$ to $(n-2)$. Hence when the average is taken for the vertices of C'_n we find that the value does not go beyond $\left\lfloor \left(\frac{n}{2}\right) \right\rfloor + 2$.

Remark 5: Let u_1, u_2, \dots, u_n be a path. The average eccentricity becomes infinite and the graph becomes disconnected for the below mentioned rotations.

$$\text{Avg Eccentricity} = \begin{cases} u_{n-1}u_n \text{ is rotated to } u_1 \\ u_1u_2 \text{ is rotated to } u_n \end{cases} = \infty$$

Remark 6: A complete binary tree of “ n ” levels requires at least one rotation to get transformed to a unicyclic graph.

Remark 7: Any tree on n vertices with $n \geq 5$ requires at least two rotations to become a unicyclic graph.

Lemma 3.3: The maximum number of rotations required for $K_{1,n}$ to be rotated to a path is $n - 2$ for any natural number n .

Proof: Leaving the last two edges or a P_3 the rest of the edges need to be rotated to form a path. Hence the number of edges that remain is $n-2$. Thus $n-2$ rotations are sufficient to transform to a path of length n .

IV. Edge Rotations on Minimum Sized Self-Centered Graphs

In this section we consider the minimal self-centered graphs of radius two proposed by Akiyama et al. [29]. We have checked these graphs for their radius and self-centeredness invariance after a single edge rotation. We also check on the average eccentricity of these graphs and give the bounds. Akiyama et al. [29] have constructed minimal self-centered graphs of radius two and shown that they one of the three classes given below.

- (i) The Petersen graph $(5,2)$.
- (ii) The graph obtained from the double star $S_{m,n}$ by adding a new vertex v and joining v to every end vertex of $S_{m,n}$.
- (iii) The graph obtained from any $K_3(a, b, c)$ by joining a new vertex w to each vertex of this $K_3(a, b, c)$.

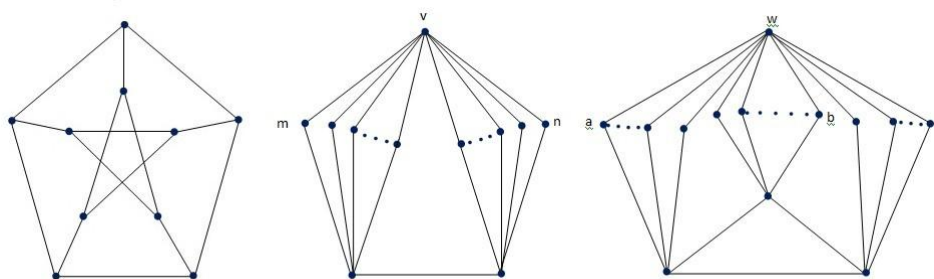


Fig 6: Minimal Self-centered graphs of radius two.

Lemma 4.1: If ‘ m ’ is the average eccentricity of any of the minimal self-centered graphs of radius two then the average eccentricity of G' is not more than $m + 1$.

Proof: For the above mentioned three graphs we find that any rotation made increases the eccentricity of the vertices. Here we observe that when the edge operation is performed on Petersen graph, the eccentricity of at most two vertices increases by one since these exist other vertices whose eccentricity may increase or remain the same., Thus resulting the change in average eccentricity by at most one. Similar argument shows that the

average eccentricity of the other two graphs increases by at most one. Thus the average eccentricity of the resulting graph does not exceed the value $m + 1$.

V. Algorithms

In the next part we develop algorithms to generate DDI graphs/ ADDI graphs from an odd and even cycle by rotating a single edge. In literature the construction of DDI graphs is considered by Itagi Huilgol et al. [21]. But constructing a DDI graph from a minimum sized DDR graph viz., cycles is an interesting problem. Hence we give algorithms to construct DDI/ADDI graphs from cycles.

Algorithm 5.1

In the following algorithm we generate DDI/ADDI graphs from an odd cycle by rotating a single edge. Let C_n be a cycle of length n .

STEP – 1: Input the cycle length n where $n \geq 5$.

STEP – 2: Find the eccentricity of the given cycle by using $ecc = \left\lfloor \left(\frac{n}{2} \right) \right\rfloor$.

STEP – 3: Consider the edge $e = uv$ to be rotated.

STEP – 4: Using the definition of rotation $G' = G - e(G) + e(\overline{G})$ perform the edge rotation.

STEP – 5: If the newly added edge results in a cycle of length less than (or equal to) the eccentricity then the graph formed is a DDI graph.

else

If $length = ecc$ or $ecc + 1$, then it is a ADDI graph .

STEP – 6 : If $length \geq ecc + 2$, then it is neither a ADDI/DDI graph.

STEP – 7 : Repeat steps 4 to 5 until $length > ecc + 2$.

STEP – 8 : STOP.

Algorithm 5.2.

In the second algorithm we develop ADDI graphs from evencycles using the concept of rotation of a single edge.

Let C_n be a cycle of length n .

STEP – 1 : Input the cycle length n where $n \geq 6$.

STEP – 2 : Find the eccentricity of the given cycle by using $e = n/2$.

STEP – 3 : Consider the edge $e = uv$ to be rotated.

STEP – 4 : Using the definition of rotation $G' = G - e(G) + e(\overline{G})$ perform the edge rotation.

STEP – 5 : If the newly added edge results in a cycle of length less than the eccentricity then the graph formed is a ADDI graph.

else

it is not a ADDI graph.

STEP – 6 : Repeat steps 4 to 5 until length of the induced cycle is less than eccentricity.

STEP – 8 : STOP.

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