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A Proof of Tarski's Fixed Point Theorem by Application of Galois Connections

**Abstract.** Two examples of Galois connections and their dual forms are considered. One of them is applied to formulate a criterion when a given subset of a complete lattice forms a complete lattice. The second, closely related to the first, is used to prove in a short way the Knaster-Tarski's fixed point theorem.

Keywords: Closure and interior operation, Galois connection, Fixed point theorem.

## 1. Introduction

For given antimonotone Galois connection defined for the complete lattices, a dual form – an appropriate monotone Galois connection (a residuated pair of mappings) is considered. The pair of closure and interior operations induced on a complete lattice by such anti- and monotone Galois connections is of our interest. Two examples of Galois connections and their dual forms are introduced in the paper. First one, considered in Sect. 3, embraces a Galois connection responsible for the dual isomorphism between a complete lattice and a closure system of subsets of a meet-generating subset of the lattice. The induced closure and interior operations are of so general form that they enable to formulate a simple criterion saying when a subset B of given complete lattice  $(A, \leq)$  forms a complete lattice with respect to the ordering  $\leq$  (Lemma 1 and Proposition 2). This criterion is applied in Sect. 4 to prove in a simple short way the Knaster-Tarski's fixed point theorem [10] (Corollary 9). The proof is constructive in the sense that it shows the explicit form of supremum and infimum of a subset in the lattice of all fixed points of a monotone mapping (cf. [2, Theorem 5.1]). This form differs from that of [2], moreover from that of [6]. The proof is also based on some simple results (inter alia Proposition 8) concerning the second example of Galois connections introduced in the paper (Sect. 4). This example is responsible

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for well-known isomorphisms between the lattice of all closure (interior) operations defined on a complete lattice  $(A, \leq)$  and the lattice of all closure (interior) systems of  $(A, \leq)$ . The induced closure and interior operations are here defined on the complete lattice of all monotone mappings of a complete lattice  $(A, \leq)$  into itself. The closure operation C induced by the antimonotone Galois connection assigns to each monotone map  $\alpha$  the least closure operation c defined on  $(A, \leq)$  such that  $\alpha \leq c$ , where  $\leq$  is the pointwise order on mappings from A to A induced by lattice ordering of  $(A, \leq)$ . In turn, the dual (monotone) Galois connection induces an interior operation I on  $(A, \leq)$  such that  $I \leq \alpha$ . A crucial point of the proof of Knaster-Tarski's theorem presented here, is the fact that the set of all fixed points of a monotone map  $\alpha$  turns out to be the intersection of the closure and interior systems of  $(A, \leq)$  corresponding to closure  $C(\alpha)$  and interior  $Int(\alpha)$  operations, respectively.

### 2. Preliminaries

The paper deals mostly with the closure and interior operations defined on a complete lattice. Given a complete lattice  $(A, \leq)$  any mapping  $C: A \longrightarrow A$ such that for each  $a \in A$ ,  $a \leq C(a)$ ,  $C(C(a)) \leq C(a)$  and C is monotone:  $a \leq b \Rightarrow C(a) \leq C(b)$ , is called a *closure operation* defined on  $(A, \leq)$ . Any subset  $B \subseteq A$  is said to be a *closure system* or *Moore family* of the lattice  $(A, \leq)$  if for each  $X \subseteq B$ ,  $\inf_A X \in B$ . Given a closure operation C on  $(A, \leq)$ , the set of all its fixed points called *closed elements*:  $\{a \in A\}$ A: a = C(a), is a closure system of  $(A, \leq)$ . Conversely, given a closure system B of  $(A, \leq)$ , the map  $C : A \longrightarrow A$  defined by  $C(a) = \inf_A \{x \in A\}$ B: a < x, is a closure operation on (A, <). The closure system B is just the set of all its closed elements. On the other hand, the closure system of all closed elements of a given closure operation C defines, in that way, just the operation C. Thus, there is a one to one correspondence between the class of all closure operations and of all closure systems of  $(A, \leq)$  (in fact it is a dual isomorphism between respective complete lattices of all closure operations and closure systems). Any closure system B of  $(A, \leq)$  forms a complete lattice with respect to the order  $\leq$  such that  $\inf_B X = \inf_A X$  and  $\sup_B X = C(\sup_A X)$ , for each  $X \subseteq B$ , where C is the closure operation corresponding to closure system B. Given a subset X of A, there exists the least closure system B of  $(A, \leq)$  such that  $X \subseteq B$ , called *generated by* X. It will be denoted here by  $[X]_{cl}$ . It is simply the intersection of all the closure

systems of  $(A, \leq)$  containing X and is of the form:  $[X]_{cl} = \{\inf_A Y : Y \subseteq X\}$ . The closure operation C corresponding to closure system  $[X]_{cl}$  is expressed by  $C(a) = \inf_A \{x \in X : a \leq x\}$ , any  $a \in A$ .

An *interior operation* and an *interior system* are the dual concepts with respect to closure ones. That is, a monotone mapping  $I: A \longrightarrow A$  such that for any  $a \in A$ ,  $I(a) \leq a, I(a) \leq I(I(a))$  is said to be an *interior* operation defined on a complete lattice  $(A, \leq)$ . Any subset B of A is called an *interior system* of the lattice  $(A, \leq)$  if for each  $X \subseteq B$ ,  $\sup_A X \in B$ . Given an interior operation I on  $(A, \leq)$  the set of all its fixed poits called *open* elements:  $\{a \in A : a = I(a)\}$ , is an interior system of  $(A, \leq)$ . Conversely, given an interior system B of  $(A, \leq)$ , the map  $I : A \longrightarrow A$  defined by  $I(a) = \sup_{A} \{x \in B : x \leq a\}$ , is an interior operation on  $(A, \leq)$ . The interior system B is just the set of all its open elements. On the other hand, the interior system of all open elements of a given interior operation I defines, in that way, just the operation I. So, as before, a similar correspondence between the class of all interior operations and interior systems, exists (which is an isomorphism of respective complete lattices of all interior operations and all interior systems of  $(A, \leq)$ ). Any interior system B of  $(A, \leq)$  forms a complete lattice with respect to the order  $\leq$  such that  $\sup_B X = \sup_A X$ and  $\inf_B X = I(\inf_A X)$ , for each  $X \subseteq B$ , where I is the interior operation corresponding to interior system B. Given a subset X of A, there exists the least interior system B of  $(A, \leq)$  such that  $X \subseteq B$ . Such an interior system is said to be *generated by* X and will be denoted as  $[X]_{in}$ . It is the intersection of all the interior systems of  $(A, \leq)$  containing X and is of the form:  $[X]_{in} = {\sup_A Y : Y \subseteq X}$ . The interior operation I corresponding to interior system  $[X]_{in}$  is defined by  $I(a) = \sup_A \{x \in X : x \leq a\}$ , any  $a \in A$ .

We shall consider the monotone and antimonotone Galois connections defined only for complete lattices. A general theory of Galois connections is to be found for example in [1, 3-5, 7].

Let us remind that while  $(A, \leq_A)$ ,  $(B, \leq_B)$  are the complete lattices, any pair of mappings  $f : A \longrightarrow B$ ,  $g : B \longrightarrow A$  such that for each  $a \in A, b \in B : b \leq_B f(a)$  iff  $a \leq_A g(b)$ , is called an *antimonotone Galois connection* for those lattices. Equivalently, such a Galois connection (f,g) fulfils the following conditions:  $a \leq_A g(f(a))$ ,  $b \leq_B f(g(b))$  for any  $a \in A, b \in B$  and f, g are antimonotone. When the pairs  $(f,g_1), (f,g_2)$ are Galois connections for the lattices  $(A, \leq_A), (B, \leq_B)$  then  $g_1 = g_2$ . The first element f of an antimonotone Galois connection (f,g) for the lattices  $(A, \leq_A), (B, \leq_B)$  is usually called a *Galois function*. A sufficient and necessary condition for a map  $f : A \longrightarrow B$  to be a Galois function is of the form:  $f(\sup_A X) = \inf_B \{f(a) : a \in X\}$ , for any  $X \subseteq A$ . Given a Galois function f, the second unique element g of the Galois connection (f,g) is given by  $g(b) = \sup_A \{a \in A : b \leq_B f(a)\}$ , for each  $b \in B$ . This mapping g satisfies the condition:  $g(\sup_B Y) = \inf_A \{g(b) : b \in Y\}$ , for any  $Y \subseteq B$ . Given a Galois connection (f,g) for the lattices  $(A, \leq_A)$ ,  $(B, \leq_B)$ , the ranges f[A], g[B] of the mappings f and g are the sets of all closed elements with respect to the closure operations  $Cl_2$ ,  $Cl_1$  respectively that are induced on B and A in the following way: for each  $a \in A$ ,  $b \in B$ ,  $Cl_2(b) = f(g(b))$ ,  $Cl_1(a) = g(f(a))$ . Since for each  $a \in g[B]$ ,  $b \in f[A] : g(f(a)) = a$ , f(g(b)) = b and moreover for any  $a_1, a_2 \in g[B] : a_1 \leq_A a_2$  iff  $f(a_2) \leq_B f(a_1)$ , so the complete lattices  $(g[B], \leq_A)$ ,  $(f[A], \leq_B)$  are dually isomorphic (with f being a dual isomorphism).

In turn, a pair  $f: A \longrightarrow B$ ,  $q: B \longrightarrow A$  such that for each  $a \in A$ ,  $b \in A$  $B: b \leq_B f(a)$  iff  $g(b) \leq_A a$ , is called a monotone Galois connection or a residuated pair of mappings for the lattices  $(A, \leq_A)$ ,  $(B, \leq_B)$ . Equivalently, a monotone Galois connection (f, g) fulfils the following conditions:  $g(f(a)) \leq_A a, b \leq_B f(g(b))$  for any  $a \in A, b \in B$  and f, g are monotone functions. When  $(f, g_1), (f, g_2)$  are residuated pairs for the lattices  $(A, \leq_A)$ ),  $(B, \leq_B)$  then  $g_1 = g_2$ . The first element f of a monotone Galois connection (f,g) for the lattices  $(A, \leq_A)$ ,  $(B, \leq_B)$  is usually called a *residuated* function while the unique second one q-a residual of f. A sufficient and necessary condition for a map  $f: A \longrightarrow B$  to be a residuated function is of the form:  $f(\inf_A X) = \inf_B \{f(a) : a \in X\}$ , for any  $X \subseteq A$ . Given a residuated function f, its residual g is expressed by  $g(b) = \inf_A \{a \in A : b \leq_B f(a)\}$ , for each  $b \in B$ . This mapping g satisfies the condition:  $g(\sup_B Y) = \sup_A \{g(b) :$  $b \in Y$ . Given a residuated pair (f, g) for the lattices  $(A, \leq_A)$ ,  $(B, \leq_B)$ , the ranges f[A], g[B] are, respectively, the sets of all closed and open elements with respect to the following closure and interior operations Cl, Int: for each  $a \in A$ ,  $b \in B$ , Cl(b) = f(q(b)), Int(a) = q(f(a)). Since for each  $a \in q[B], b \in f[A]$ : q(f(a)) = a, f(q(b)) = b and moreover for any  $a_1, a_2 \in g[B]$ :  $a_1 \leq_A a_2$  iff  $f(a_1) \leq_B f(a_2)$ , so the complete lattices  $(g[B], \leq_A), (f[A], \leq_B)$  are isomorphic (with f being an isomorphism).

From the very definition of Galois connections it follows that any antimonotone Galois connection (f,g) for the lattices  $(A, \leq_A)$ ,  $(B, \leq_B)$  is simultaneously a residuated pair for the lattices  $(A, \leq_A^{\sim})$ ,  $(B, \leq_B)$ , where  $\leq_A^{\sim}$  is the converse ordering to  $\leq_A$ . Taking this into account, having defined a Galois function  $f_{\leq_A} : A \longrightarrow B$  for the complete lattices  $(A, \leq_A)$ ,  $(B, \leq_B)$ (we write down the parameter:  $\leq_A$ , on which the function may depend as an essential one, however in general there are the other parameters which may occur in a definition of Galois function) let us consider a mapping  $f_{\leq_{\widetilde{A}}}: A \longrightarrow B$  which is defined exactly in the same way as the function  $f_{\leq_A}$  except that instead of the parameter  $\leq_A$  the converse relation is applied. Notice that when  $f_{\leq_A}$  being a Galois function fulfils the condition:  $f_{\leq_A}(\sup_{\leq_A} X) = \inf_{\leq_B} \{f_{\leq_A}(a) : a \in X\}$ , the mapping  $f_{\leq_{\widetilde{A}}}$  has to satisfy the following one:  $f_{\leq_{\widetilde{A}}}(\inf_{\leq_A} X) = \inf_{\leq_B} \{f_{\leq_{\widetilde{A}}}(a) : a \in X\}$ , any  $X \subseteq A$ , that is,  $f_{\leq_{\widetilde{A}}}$  is a residuated function for the lattices  $(A, \leq_A)$ ,  $(B, \leq_B)$ . Let us call such a residuated function the *dual residuated function with respect* to  $f_{\leq_A}$ . Moreover, when (f, g) is an antimonotone Galois connection let us call the residuated pair  $(f_d, g_d)$ , where  $f_d$  is the dual residuation function with respect to f, the *dual residuated pair* (or the *dual Galois connection*) with respect to (f, g). Obviously, one can start not from a Galois but a residuated function (residuated pair) and define the dual Galois function (the dual antimonotone Galois connection).

Having at our disposal the Galois connections: (f, g),  $(f_d, g_d)$  for the complete lattices  $(A, \leq_A)$ ,  $(B, \leq_B)$  we are especially interested in the interiorclosure pair (Int, C) of operations on  $(A, \leq_A)$ , where  $Int = f_d \circ g_d$  and  $C = f \circ g$  (the closure operation C was denoted by  $Cl_1$  above).

In the sequel we consider two important examples of antimonotone Galois connections and their dual forms. First one enables to formulate a simple criterion saying when a given subset of a complete lattice forms a complete lattice. The second example, closely related to the first, has rather unexpected applications. It enables a very simple proving of the Knaster-Tarski's fixed point theorem.

### 3. A Criterion of Being a Complete Lattice

Let  $(A, \leq)$  be any complete lattice and  $B \subseteq A$ . The following pair of mappings:  $f: A \longrightarrow \wp(B), g: \wp(B) \longrightarrow A$  defined by  $f(a) = \{x \in B : a \leq x\}$ , any  $a \in A$  and  $g(X) = \inf_A X$ , any  $X \subseteq B$ , forms an antimonotone Galois connection for the lattices  $(A, \leq), (\wp(B), \subseteq)$ . The dual residuated function with respect to f is then of the form:  $f_d(a) = \{x \in B : x \leq a\}$  and its residual is defined by  $g_d(X) = \inf_A \{a \in A : X \subseteq f_d(a)\} = \inf_A \{a \in A : X \subseteq \{x \in B : x \leq a\}\} = \sup_A X$ , as one could expect.

These Galois connections are responsible for well-known isomorphisms of a complete lattice and a lattice of subsets of a given meet- or join-generating subset of the lattice. A subset B of a complete lattice  $(A, \leq)$  is said to be *join*generating (meet-generating, cf. for example [5]) or *join-dense* (meet-dense, e.g. [8]) iff for each  $a \in A$ , there is an  $X \subseteq B$  such that  $a = \sup_A X$  (a =  $\inf_A X$ ). For example, the set of all compact elements of an algebraic lattice is just its join-generating subset.

It is clear that the restriction of the map f to the set  $\{\inf_A X : X \subseteq B\}$ (which is the closure system generated by B) is a dual isomorphism of the lattice  $(\{\inf_A X : X \subseteq B\}, \leq)$  of all closed elements with respect to the closure operation  $C_B = f \circ g$  to the lattice  $(\{B \cap [a) : a \in A\}, \subseteq)$  (which is the closure system of  $(\wp(B), \subseteq)$  corresponding to closure operation  $g \circ f$ ; here  $[a) = \{x \in A : a \leq x\}$ ). Similarly, the restriction of the map  $f_d$ to the set  $\{\sup_A X : X \subseteq B\}$  (which is the interior system generated by B) is an isomorphism of the lattice  $(\{\sup_A X : X \subseteq B\}, \leq)$  of all open elements with respect to the interior operation  $I_B = f_d \circ g_d$  to the lattice  $(\{B \cap (a] : a \in A\}, \subseteq)$  (being the closure system of  $(\wp(B), \subseteq)$  corresponding to closure operation  $g_d \circ f_d$ ; here  $(a] = \{x \in A : x \leq a\}$ ).

One can easily see from their definitions that the operations  $I_B, C_B$  are of the following general form, for any  $a \in A$ :

(1) 
$$I_B(a) = \sup_A \{ x \in B : x \le a \},\$$

(2) 
$$C_B(a) = \inf_A \{x \in B : a \le x\}.$$

They simply corrrespond to the interior and to closure systems of  $(A, \leq)$ generated by B, respectively. The pair  $(I_B, C_B)$  is a generalization of the notion of so-called pair of interior-closure operations associated on a given subset of a complete lattice, introduced in [9] and widely applied there. In case a subset B forms a complete sublattice of the lattice  $(A, \leq)$ , the pair  $(I_B, C_B)$  becomes just an interior-closure pair of operations associated on B. The existence of an interior-closure pair of operations associated on B is a necessary and sufficient condition for  $(B, \leq)$  to be a complete sublattice of  $(A, \leq)$  (cf. [9]). This criterion will be now generalized in order to provide the sufficient and necessary conditions for the poset  $(B, \leq)$  to be a complete lattice. Let us start from the crucial lemma.

LEMMA 1. Let  $D, O \subseteq A$  be any closure and interior systems of a complete lattice  $(A, \leq)$ , respectively. Then the following conditions are equivalent:

- (i) for each  $a \in O$ ,  $C_D(a) \in O$ ,
- (ii) for each  $a \in A$ ,  $C_D(I_O(a)) \in O$ ,
- (iii) for each  $a \in A$ ,  $I_O(C_D(a)) \in D$ ,
- (iv) for each  $a \in D$ ,  $I_O(a) \in D$ ,

where the operations  $I_O, C_D$  are defined by (1) and (2), respectively, for the sets O, D instead of B. Moreover, any of these conditions implies that the

poset  $(D \cap O, \leq)$  is a complete lattice in which for any  $X \subseteq D \cap O$ ,  $\sup X = C_D(\sup_A X)$  and  $\inf X = I_O(\inf_A X)$ . The inverse implication in general does not hold.

PROOF. Suppose that the subsets D and O of A are closure and interior systems of a complete lattice  $(A, \leq)$ , respectively. The equivalences  $(i) \Leftrightarrow$  $(ii), (iii) \Leftrightarrow (iv)$  are obvious. In order to show the implication  $(ii) \Rightarrow (iii)$ assume that for each  $a \in A$ ,  $C_D(I_O(a)) = I_O(C_D(I_O(a)))$ . Then given  $a \in A$ we have  $C_D(I_O(C_D(a))) = I_O(C_D(I_O(C_D(a))))$ . Since  $I_O(C_D(a)) \leq C_D(a)$ so  $C_D(I_O(C_D(a))) \leq C_D(a)$  ( $C_D$  is monotone and idempotent). Therefore,  $I_O(C_D(I_O(C_D(a)))) \leq I_O(C_D(a))$  (by monotonicity of  $I_O$ ) which together with the last identity implies that  $C_D(I_O(C_D(a))) \leq I_O(C_D(a))$  so we obtain (iii). The proof from (iii) to (ii) goes analogously (by dual argument).

In order to prove the second part of lemma suppose (i) and consider an  $X \subseteq D \cap O$ . Then since O is an interior system we have  $\sup_A X \in O$ . So from (i) it follows that  $C_D(\sup_A X) \in D \cap O$ . Now, given any  $a \in X$  we have  $a \leq \sup_A X \leq C_D(\sup_A X)$ , so  $C_D(\sup_A X)$  is an upper bound of X in the poset  $(D \cap O, \leq)$ . When  $z \in D \cap O$  is such an upper bound we obtain:  $\sup_A X \leq z$ , therefore  $C_D(\sup_A X) \leq C_D(z) = z$ . In this way,  $C_D(\sup_A X)$  is the least upper bound of X in  $(D \cap O, \leq)$ . The form of inf X in this poset follows from the condition (iv) in a similar way.

Finally, in order to show that none of the conditions (i) - (iv) needs to be true when a poset  $(D \cap O, \leq)$  is a complete lattice, take for example a 4-element chain: 0 < a < b < 1 and consider  $D = \{0, b, 1\}, O = \{0, a, 1\}$ .

Now let us formulate our criterion saying when a subset of given complete lattice  $(A, \leq)$  forms a complete lattice with respect to the order  $\leq$ .

PROPOSITION 2. Let  $(A, \leq)$  be a complete lattice and  $B \subseteq A$ . Consider the operations  $I_B, C_B$  defined by (1), (2). The following conditions are equivalent:

- (a) for each  $a \in A$ ,  $C_B(I_B(a)) \in B$ ,
- (b) for each  $a \in A$ ,  $I_B(C_B(a)) \in B$ ,
- (c)  $(B, \leq)$  is a complete lattice such that for any  $X \subseteq B$ ,  $\sup X = C_B(\sup_A X)$  and  $\inf X = I_B(\inf_A X)$ .

PROOF. Let  $B \subseteq A$ . Put  $D = [B]_{cl}$ ,  $O = [B]_{in}$ . Then we have immediately  $B \subseteq D \cap O$  and  $C_D = C_B$ ,  $I_O = I_B$ .

 $(a) \Rightarrow (b)$  & (c): Assume that (a) holds. Then the condition (ii) of Lemma 1 is satisfied. Moreover, taking any  $a \in D \cap O$  we have

 $C_B(I_B(a)) = a$  so from (a) it follows that  $a \in B$ , consequently,  $B = D \cap O$ . Thus, on one hand, from (ii) and Lemma 1 it follows that (iii) of Lemma 1 holds which leads to (b). On the other hand, simultaneously from (ii) and Lemma 1 it follows that (c) holds true.

 $(b) \Rightarrow (a)$ : By the dual argument with respect to the proof of implication  $(a) \Rightarrow (b)$ .

 $(c) \Rightarrow (a)$ : Suppose that (c) holds. Let  $a \in A$ . Since  $I_B(a) \in [B]_{in}$  so  $I_B(a) = \sup_A X$  for some  $X \subseteq B$ . Therefore,  $C_B(I_B(a)) = C_B(\sup_A X) = \sup X$  by (c). Thus,  $C_B(I_B(a)) \in B$ .

# 4. The Galois Connections Involving Monotone Mappings on Complete Lattices

Let  $(A, \leq)$  be a complete lattice and *Mon*-the class of all monotone mappings from A to A. Obviously, the poset  $(Mon, \leq)$  is a complete sublattice of the complete lattice  $(A^A, \leq)$  of all the mappings from A to A, where for any  $\alpha, \beta \in A^A$ ,  $\alpha \leq \beta$  iff for all  $x \in A$ ,  $\alpha(x) \leq \beta(x)$ . For any  $F \subseteq Mon$ ,  $(\sup F)(a) = \sup_A \{\alpha(a) : \alpha \in F\}$  and  $(\inf F)(a) = \inf_A \{\alpha(a) : \alpha \in F\}$ , for each  $a \in A$ .

The main goal of this section is to prove the Knaster-Tarski's fixed point theorem using a special Galois connection. This Galois connection turns out to be significant also from the other point of view. It is responsible for well-known dual isomorphism between the complete lattice of all closure operations defined on the complete lattice  $(A, \leq)$  and the complete lattice of all closure systems of  $(A, \leq)$ . The connection is of the form: f : (Mon, $\leq) \longrightarrow (\wp(A), \subseteq)$  is a mapping defined by  $f(\alpha) = \{x \in A : \alpha(x) \leq x\}$  and  $g : (\wp(A), \subseteq) \longrightarrow (Mon, \leq)$  is such that for any  $B \subseteq A, g(B) : A \longrightarrow A$  is defined by  $g(B)(a) = \inf_A \{x \in B : a \leq x\} = \inf_A (B \cap [a])$ . It is obvious that g(B) for each  $B \subseteq A$  is monotone. Notice simply that given  $B \subseteq A, g(B)$ is just the closure operation  $C_B$  from the previous section.

LEMMA 3. (f,g) is a Galois connection, i.e., f,g are antimonotone, for each  $\alpha \in Mon, \ \alpha \leq g(f(\alpha))$  and for any  $B \subseteq A, \ B \subseteq f(g(B))$ .

PROOF. The proof that both f, g are antimonotone is straightforward. In order to show that given  $\alpha \in Mon$ ,  $\alpha \leq g(f(\alpha))$ , notice that given  $a \in$  $A, g(f(\alpha))(a) = \inf_A \{x \in A : \alpha(x) \leq x \& a \leq x\}$ . Consider any  $x \in A$ such that  $\alpha(x) \leq x$  and  $a \leq x$ . Then since the map  $\alpha$  is monotone we have:  $\alpha(a) \leq \alpha(x)$  which implies that  $\alpha(a) \leq x$ . This means that  $\alpha(a)$  is a lower bound of the set  $\{x \in A : \alpha(x) \leq x \& a \leq x\}$  in the lattice  $(A, \leq)$ . Therefore,  $\alpha(a) \leq \inf_A \{x \in A : \alpha(x) \leq x \& a \leq x\}$ , that is  $\alpha(a) \leq g(f(\alpha))(a)$ . To the end, in order to prove that for all  $B \subseteq A$ ,  $B \subseteq f(g(B))$  take any  $a \in B$ . Our goal is to show that  $g(B)(a) \leq a$ . However, in case  $a \in B$  we have:  $\inf_A \{x \in B : a \leq x\} = a$ , so g(B)(a) = a.

Now, consider the closure operations induced by the Galois connection  $(f,g), Cl_1 : Mon \longrightarrow Mon$  and  $Cl_2 : \wp(A) \longrightarrow \wp(A)$  defined by  $Cl_1(\alpha) = g(f(\alpha))$ , for any  $\alpha \in Mon$  and  $Cl_2(B) = f(g(B))$ , for each  $B \subseteq A$ . Obviously,  $\{\alpha \in Mon : Cl_1(\alpha) = \alpha\} = g[\wp(A)]$  and  $\{B \subseteq A : Cl_2(B) = B\} = f[Mon]$ . Moreover, the mapping f restricted to the set  $\{\alpha \in Mon : Cl_1(\alpha) = \alpha\}$  is a dual isomorphism between the posets  $(\{\alpha \in Mon : Cl_1(\alpha) = \alpha\}, \leq), (\{B \subseteq A : Cl_2(B) = B\}, \subseteq).$ 

One may characterize the sets of all closed elements with respect to the first and to the second closure operations in the following way.

- PROPOSITION 4. (1) For any  $\alpha \in Mon$ ,  $Cl_1(\alpha) = \alpha$  iff  $\alpha$  is a closure operation on  $(A, \leq)$ .
  - (2) For any  $B \subseteq A$ ,  $Cl_2(B) = B$  iff for any  $X \subseteq B$ ,  $\inf_A X \in B$ , that is B is a closure system of the lattice  $(A, \leq)$ .

PROOF. For (1) ( $\Rightarrow$ ): Assume that  $Cl_1(\alpha) = \alpha$ . Then  $\alpha = g(B)$  for some  $B \subseteq A$ . So  $\alpha$  is the closure operation  $C_B$  on  $(A, \leq)$  from the previous section.

( $\Leftarrow$ ): Assume that  $\alpha$  is a closure operation on  $(A, \leq)$ . Our goal is to show that  $g(f(\alpha)) \leq \alpha$ . For each  $a \in A$  we have  $g(f(\alpha))(a) = \inf_A \{x \in A : \alpha(x) \leq x \& a \leq x\}$ . From the assumption it follows that given  $a \in A$ ,  $\alpha(\alpha(a)) \leq \alpha(a)$  and  $a \leq \alpha(a)$ , so  $\alpha(a) \in \{x \in A : \alpha(x) \leq x \& a \leq x\}$ , thus  $\inf_A \{x \in A : \alpha(x) \leq x \& a \leq x\} \leq \alpha(a)$ , that is  $g(f(\alpha))(a) \leq \alpha(a)$ .

For (2) ( $\Rightarrow$ ): Assume that  $Cl_2(B) = B$  and  $X \subseteq B$ . Then obviously,  $B = f(\alpha)$  for some  $\alpha \in Mon$ , that is,  $B = \{x \in A : \alpha(x) \leq x\}$  for some  $\alpha \in Mon$ . So we have furthermore  $X \subseteq \{x \in A : \alpha(x) \leq x\}$ . Hence, taking any  $a \in X$  into account we have  $\alpha(a) \leq a$  while from the monotonicity of  $\alpha$  it follows that  $\alpha(\inf_A X) \leq \alpha(a)$  (for  $\inf_A X \leq a$ ). Thus  $\alpha(\inf_A X) \leq a$ , so  $\alpha(\inf_A X)$  is a lower bound of X, therefore,  $\alpha(\inf_A X) \leq \inf_A X$ . This means that  $\inf_A X \in B$ .

( $\Leftarrow$ ): Assume that for all  $X \subseteq B$ ,  $\inf_A X \in B$ . It is sufficient to show that  $f(g(B)) \subseteq B$ . We have  $f(g(B)) = \{a \in A : g(B)(a) \le a\} = \{a \in A : \inf_A \{x \in B : a \le x\} \le a\} = \{a \in A : \inf_A \{x \in B : a \le x\} = a\} = \{a \in A : \inf_A (B \cap [a)) = a\}$ . So let  $a \in f(g(B))$ . Then  $\inf_A (B \cap [a)) = a$ . Since  $B \cap [a] \subseteq B$  so from the assumption it follows that  $\inf_A (B \cap [a)) \in B$ , that is,  $a \in B$ . As one may see, Proposition 4 yields the above-mentioned correspondence between the closure operations and closure systems of given complete lattice.

- COROLLARY 5. (1) For any monotone mapping  $\alpha : A \longrightarrow A$ ,  $Cl_1(\alpha)$  is the least closure operation  $c : A \longrightarrow A$  such that  $\alpha \leq c$ . Explicitly, for any  $a \in A : Cl_1(\alpha)(a) = C_{f(\alpha)}(a) = \inf_A \{x \in A : \alpha(x) \leq x \& a \leq x\}.$
- (2) For any  $B \subseteq A$ ,  $Cl_2(B)$  is the least closure system  $Z \subseteq A$  such that  $B \subseteq Z$  (i.e.  $Cl_2(B) = [B]_{cl}$ ). Explicitly,  $Cl_2(B) = \{\inf_A X : X \subseteq B\}$ .

PROOF. It is obvious that given any poset  $(Y, \leq)$  and a closure operation  $Cl: Y \longrightarrow Y$ , for any  $y \in Y$ , Cl(y) is the least element  $y' \in \{x \in Y : x = Cl(x)\}$  such that  $y \leq y'$ . So we obtain the first statements of (1) and (2) due to Proposition 4 since  $\{\alpha \in Mon : Cl_1(\alpha) = \alpha\}$  is the class of all the closure operations mapping A into A, and  $\{B \subseteq A : Cl_2(B) = B\}$  is the family of all the closure systems contained in A. The explicit form of the operation  $Cl_1$  immediately follows from its definition (comp. the proof for (1) ( $\Leftarrow$ ) of Proposition 4). In order to show the explicit form of  $Cl_2$  we have to show, according to the proof for (2) ( $\Leftarrow$ ) of Proposition 4, that  $\{a \in A : \inf_A(B \cap [a)) = a\} = \{\inf_A X : X \subseteq B\}$ . The inclusion ( $\subseteq$ ) is obvious. In order to prove the inverse inclusion take any  $X \subseteq B$ . Then  $X \subseteq \{x \in B : \inf_A X \leq x\} = B \cap [\inf_A X)$ . Hence  $\inf_A(B \cap [\inf_A X)) \leq \inf_A X$ . However, on the other hand, the element  $\inf_A X$  is a lower bound of the set  $B \cap [\inf_A X)$ . So  $\inf_A X \leq \inf_A (B \cap [\inf_A X))$  and finally  $\inf_A X = \inf_A (B \cap [\inf_A X))$ . Thus,  $\inf_A X \in \{a \in A : \inf_A (B \cap [a)) = a\}$ .

Now let us consider the dual residuated pair of mappings with respect to Galois connection (f, g). The dual residuated function  $f_d$  should be defined by changing in the definition of f the order  $\leq$  defined on Mon into its inverse order. But the order in the complete lattice of all monotone mappings from A to A is in turn defined by the order of the lattice  $(A, \leq)$ . So taking the inverse order on mappings means to take into consideration the inverse order of  $\leq$  on A. Therefore we put  $f_d(\alpha) = \{x \in A : x \leq \alpha(x)\}$ . One can check that so defined map fulfils the condition for being a residuated function for the complete lattices  $(Mon, \leq), (\wp(A), \subseteq)$ : given  $F \subseteq Mon, f_d(\inf F) = \{x \in A : x \leq (\inf F)(x)\} = \{x \in A : x \leq \inf_A \{\alpha(x) : \alpha \in F\}\} = \bigcap\{\{x \in A : x \leq \alpha(x)\}: \alpha \in F\} = \bigcap\{f_d(\alpha) : \alpha \in F\}$ .

According to the general definition of a residual, we have for any  $B \subseteq A$ :  $g_d(B) = \inf\{\alpha \in Mon : B \subseteq f_d(\alpha)\}$ . So for each  $a \in A$ ,  $g_d(B)(a) = \inf_A\{\alpha(a) : \alpha \in Mon \& B \subseteq f_d(\alpha)\} = \inf_A\{\alpha(a) : \alpha \in Mon \& B \subseteq \{x \in A : x \leq \alpha(x)\}\}$ . It is easily seen that given  $a \in A$ ,  $g_d(B)(a)$  is an upper bound of

the set  $\{x \in B : x \leq a\}$  in the lattice  $(A, \leq)$ . On the other hand, consider any upper bound z of the set  $\{x \in B : x \leq a\}$ , that is,  $\forall x \in B \ (x \leq a \Rightarrow x \leq z)$ . Then a monotone mapping  $\alpha_z$  defined on A by  $\alpha_z(x) = z$  whenever  $x \leq a$ otherwise  $\alpha_z(x) = 1_A$  (the unit of the complete lattice  $(A, \leq)$ ), is such that  $B \subseteq \{x \in A : x \leq \alpha_z(x)\}$ . From this and the fact:  $\alpha_z(a) = z$ , it follows that  $z \in \{\alpha(a) : \alpha \in Mon \& B \subseteq \{x \in A : x \leq \alpha(x)\}\}$  and consequently,  $g_d(B)(a) \leq z$ . Finally,  $g_d(B)(a) = \sup_A \{x \in B : x \leq a\}$ . So, given  $B \subseteq A$ , the mapping  $g_d(B)$  is just the interior operation  $I_B$  from the previous section so it is monotone.

Now, one can consider the interior operation  $Int : Mon \longrightarrow Mon$  and the closure operation  $Cl : \wp(A) \longrightarrow \wp(A)$  induced by the residuated pair  $(f_d, g_d)$  that is defined by  $Int(\alpha) = g_d(f_d(\alpha))$ , for any  $\alpha \in Mon$  and Cl(B) = $f_d(g_d(B))$ , for each  $B \subseteq A$ . Furthermore, firstly,  $\{\alpha \in Mon : Int(\alpha) = \alpha\} =$  $g_d[\wp(A)]$  and  $\{B \subseteq A : Cl(B) = B\} = f_d[Mon]$ . Secondly, the mapping  $f_d$ restricted to the set  $\{\alpha \in Mon : Int(\alpha) = \alpha\}$  is an isomorphism between the posets ( $\{\alpha \in Mon : Int(\alpha) = \alpha\}, \leq$ ), ( $\{B \subseteq A : Cl(B) = B\}, \subseteq$ ). Thus, the following proposition is responsible for an isomorphism between the complete lattices of all interior operations and interior systems defined on given complete lattice.

- PROPOSITION 6. (1) For any  $\alpha \in Mon$ ,  $Int(\alpha) = \alpha$  iff  $\alpha$  is an interior operation on  $(A, \leq)$ .
  - (2) For any  $B \subseteq A$ , Cl(B) = B iff for each  $Y \subseteq B$ ,  $\sup_A Y \in B$ , that is B is an interior system in the lattice  $(A, \leq)$ .

**PROOF.** Analogous to the proof of Proposition 4, by dual argument.

- COROLLARY 7. (1) For any monotone mapping  $\alpha : A \longrightarrow A$ ,  $Int(\alpha)$  is the greatest interior operation  $I : A \longrightarrow A$  such that  $I \leq \alpha$ . Explicitly, for any  $a \in A$ :  $Int(\alpha)(a) = I_{f_d(\alpha)}(a) = \sup_A \{x \in A : x \leq \alpha(x) \& x \leq a\}.$
- (2) For any  $B \subseteq A$ , Cl(B) is the least interior system  $Z \subseteq A$  such that  $B \subseteq Z$  (that is  $Cl(B) = [B]_{in}$ ). Explicitly,  $Cl(B) = \{\sup_A Y : Y \subseteq B\}$ .

PROOF. Analogous to the proof of Corollary 5, by dual argument.

EXAMPLE. Consider the lattice  $(Mon, \leq)$  of all monotone mappings defined on 4-element lattice  $(\{0, a, b, 1\}, \leq)$  with a, b – incomparable elements (Fig. 1). There are 7 closure and 7 interior operations in *Mon* (Fig. 2). The set *Mon* may be divided into 7 equivalent classes modulo the equivalence relation  $\theta_f$  induced on *Mon* by f (that is  $\alpha \theta_f \beta$  iff  $f(\alpha) = f(\beta)$ ) as well as by  $f_d$ .



Figure 1. The lattice  $(Mon, \leq)$  for the lattice  $(\{0, a, b, 1\}, \leq)$ 

Here we write down the explicit form of each equivalence class modulo  $\theta_f$ : {1111, a111, a1a1, b111, bb11}, {aa11, aaa1, aaaa}, {0111, 01a1, 0b11, 0ba1}, {b1b1, bbb1, bbbb}, {0a11, 0aa1, 0011, 0aaa, 00a1, 00aa}, {01b1, 0101, 0bb1, 0b01, 0bbb, 0b0b}, {0ab1, 0a01, 00b1, 0a0a, 0001, 00bb, 000a, 000b, 0000}. In each class at the first place a closure operation occurs. This is the unique



Figure 2. The lattice of all closure and interior operations on the lattice  $(\{0,a,b,1\},\leq)$ 

closure operation in a given class, being the greatest element of it, denoted so far as  $Cl_1(\alpha)$  or  $C_{f(\alpha)}$  for any map  $\alpha$  from the equivalence class.

Now let us proceed to a proof of Knaster-Tarski's theorem. To this aim first let us remind that given a monotone mapping  $\alpha : A \longrightarrow A$  we have:  $Cl_1(\alpha) = g(f(\alpha)) = C_{f(\alpha)}$ . Explicitly, for each  $a \in A$ ,  $C_{f(\alpha)}(a) = \inf_A \{x \in f(\alpha) : a \leq x\} = \inf_A \{x \in A : \alpha(x) \leq x \& a \leq x\}$ . The set  $f(\alpha) = \{x \in A : \alpha(x) \leq x\}$ , is the closure system corresponding (by dual isomorphism g) to closure operation  $C_{f(\alpha)}$ , so  $f(C_{f(\alpha)}) = \{x \in A : C_{f(\alpha)}(x) \leq x\} = \{x \in A : C_{f(\alpha)}(x) = x\} = f(\alpha)$ . Moreover,  $Int(\alpha) = g_d(f_d(\alpha)) = I_{f_d(\alpha)}$ , that is 
$$\begin{split} I_{f_d(\alpha)}(a) &= \sup_A \{ x \in f_d(\alpha) : x \leq a \} = \sup_A \{ x \in A : x \leq \alpha(x) \& x \leq a \}. \\ \text{The set } f_d(\alpha) &= \{ x \in A : x \leq \alpha(x) \}, \text{ is the interior system corresponding} \\ (\text{by isomorphism } g_d) \text{ to interior operation } I_{f_d(\alpha)}, \text{ so } f_d(I_{f_d(\alpha)}) &= \{ x \in A : x \leq I_{f_d(\alpha)}(x) \} \\ &= \{ x \in A : I_{f_d(\alpha)}(x) = x \} = f_d(\alpha). \text{ Since } \alpha \text{ is monotone,} \\ \text{both systems: } f(\alpha), f_d(\alpha) \text{ are closed on } \alpha \text{ conceived as an unary operation} \\ \text{on } A. \end{split}$$

PROPOSITION 8. For all  $\alpha \in Mon$ :

- (1) the interior system  $f_d(\alpha)$  is closed on the operation  $C_{f(\alpha)}$ : for any  $a \in f_d(\alpha), C_{f(\alpha)}(a) \in f_d(\alpha),$
- (2) the closure system  $f(\alpha)$  is closed on the operation  $I_{f_d(\alpha)}$ : for any  $a \in f(\alpha)$ ,  $I_{f_d(\alpha)}(a) \in f(\alpha)$ .

PROOF. Assume that  $\alpha : A \longrightarrow A$  is any monotone mapping. In order to show (1) suppose that  $a \in f_d(\alpha)$ . Hence and from the assumption it follows that  $a \leq \alpha(a) \leq \alpha(C_{f(\alpha)}(a))$ . Moreover,  $\alpha(C_{f(\alpha)}(a)) \in f(\alpha)$  for  $C_{f(\alpha)}(a)$  is a closed element and the set  $f(\alpha)$  of all closed elements with respect to  $C_{f(\alpha)}$ is closed on  $\alpha$ . In this way,  $\alpha(C_{f(\alpha)}(a)) \in \{x \in f(\alpha) : a \leq x\}$ . Therefore,  $\inf_A \{x \in f(\alpha) : a \leq x\} \leq \alpha(C_{f(\alpha)}(a))$ , that is,  $C_{f(\alpha)}(a) \leq \alpha(C_{f(\alpha)}(a))$ . This means that  $C_{f(\alpha)}(a) \in f_d(\alpha)$ . Analogously for (2). Obviously, the conditions (1), (2) are equivalent due to Lemma 1 (i)  $\Leftrightarrow$  (iv).

COROLLARY 9. (The Knaster-Tarski's fixed point theorem [10]) Given a complete lattice  $(A, \leq)$  and a monotone function  $\alpha : A \longrightarrow A$ , the poset  $(B, \leq)$ , where  $B = \{x \in A : x = \alpha(x)\}$ , is a complete lattice in which for any  $X \subseteq B$ ,  $\sup X = C_{f(\alpha)}(\sup_A X)$  and  $\inf X = I_{f_d(\alpha)}(\inf_A X)$ .

PROOF. By simple application of Lemma 1 for  $D = f(\alpha)$ ,  $O = f_d(\alpha)$ . Any of conditions (i) - (iv) of Lemma 1 is satisfied due to Proposition 8.

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