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# MINIMUM DEGREE CONDITIONS FOR TILINGS IN GRAPHS AND HYPERGRAPHS

by

ANDREW LIGHTCAP

Under the Direction of Dr. Yi Zhao

## ABSTRACT

We consider tiling problems for graphs and hypergraphs. For two graphs  $G$  and  $F$ , an  $F$ -tiling of  $G$  is a subgraph of  $G$  consisting of only vertex disjoint copies of  $F$ . By using the absorbing method we give a short proof that in a balanced tripartite graph  $G$ , if every vertex is adjacent to  $(2/3 + \gamma)$  of the vertices in each of the other vertex partitions, then  $G$  has a  $K_3$  tiling. Previously Magyar and Martin [14] proved the same result (without  $\gamma$ ) by using the Regularity Lemma.

In a 3-uniform hypergraph  $\mathcal{H}$ , let  $\delta_2(\mathcal{H})$  denote the minimum number of edges that contain  $\{u, v\}$  for all pairs  $\{u, v\}$  of vertices. We show that if  $\delta_2(\mathcal{H}) \geq \left(1 - \frac{2}{k(k-2)}\right)n$  there exists a  $K_k^3$ -tiling of  $\mathcal{H}$  that misses at most  $k^2$  vertices of  $\mathcal{H}$ . On the other hand, we show that there exist hypergraphs  $\mathcal{H}$  such that  $\delta_2(\mathcal{H}) = \left(1 - \frac{1}{k}\right)n - 2$  and  $\mathcal{H}$  does not have a perfect  $K_k^3$ -tiling. These extend the results of Pikhurko [17] on  $K_4^3$ -tilings.

INDEX WORDS: Graph tiling, Graph packing, Absorbing method, Hypergraph Codegree

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by

ANDREW LIGHTCAP

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MINIMUM DEGREE CONDITIONS FOR TILINGS IN GRAPHS AND HYPERGRAPHS

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This thesis is dedicated to Luzy, who stands beside me on all of my adventures.

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## Chapter 1

### INTRODUCTION

For two graphs  $G$  and  $F$ , an  $F$ -tiling (or  $F$ -packing) of  $G$  is a subgraph of  $G$  consisting of vertex disjoint copies of  $F$ . When  $F$  is a single (hyper)edge we call an  $F$ -tiling a matching. If the  $F$ -tiling covers all of the vertices of  $G$  we say that the tiling is perfect or refer to the tiling as an  $F$ -factor. For a perfect tiling to exist the order of  $F$  must divide the order of  $G$ .

The purpose of this paper is to determine bounds on the minimum degree necessary to ensure a perfect or near perfect  $F$ -tiling. An early result by Dirac [6] proves that any graph on  $n$  vertices with minimum degree at least  $n/2$  is Hamiltonian. This result allows us to obtain a perfect matching in  $G$  by deleting every other edge from the Hamiltonian cycle. For  $F = K_h$ , the complete graph on  $h$  vertices, Hajnal and Szemerédi [8] provide the following result: *If  $G$  is a graph with  $hk$  vertices and minimum degree at least  $(h - 1)k$ , then  $G$  contains  $k$  vertex disjoint copies of  $K_h$ .* Later, using Szemerédi's Regularity Lemma [22], Alon and Yuster [2, 3] were able to provide minimum degree conditions that guarantee an  $F$ -factor for arbitrary  $F$ . Kühn and Osthus [12] were able to find the best possible minimum degree conditions for finding an  $F$ -factor.

Tiling in multipartite graphs has a shorter history. A graph  $G$  is called  $r$ -partite if the vertex set  $V(G)$  can be partitioned in  $r$  sets  $V_1, \dots, V_r$  such any that two vertices  $u, v \in V_i$  are not adjacent. The Marriage Theorem by König and Hall (see e.g. [4]) implies that a bipartite graph ( $r = 2$ )  $G$  with partition sets of size  $n$  contains a 1-factor if  $\delta(G) \geq n/2$ . In an  $r$ -partite graph  $G$  with  $r \geq 2$ , let  $\bar{\delta}(G)$  be the minimum degree from a vertex in one partition set to each other partition set (so  $\bar{\delta}(G) = \delta(G)$  when  $r = 2$ ). An  $r$ -partite graph is balanced if all partition sets have the same order.

Fischer [7] conjectured the following  $r$ -partite version of the Hajnal-Szemerédi Theorem and

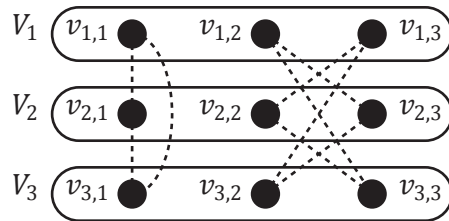


Figure 1.1. Representation of  $\Gamma_3$  with dotted lines corresponding to non-edges

proved it asymptotically for  $r = 3, 4$ : if  $G$  is an  $r$ -partite graph with  $n$  vertices in each partition set and  $\bar{\delta}(G) \geq \frac{r-1}{r}n$ , then  $G$  contains a  $K_r$ -factor. Magyar and Martin [14] used the following theorem to show that Fischer's conjecture is slightly wrong for  $r = 3$  (off by only 1): For  $G$  a balanced tripartite graph on  $3N$  vertices with  $\bar{\delta}(G) \geq (2/3)N + 1$  then  $G$  contains a perfect  $K_3$ -tiling. As written, this is a weaker form of the actual theorem, as they prove that  $G$  can be perfectly tiled with triangles when  $\bar{\delta}(G) \geq (2/3)N$  as long as it is not the graph  $\Gamma_3(N/3)$ . The case when  $G$  is  $\Gamma_3(N/3)$  is what disproves Fischer's conjecture and necessitates the extra edge to complete the tiling. Notice in Figure 1 that there can be no  $K_3$ -tiling of  $\Gamma_3$ . To form  $\Gamma_3(N/3)$ , replace each vertex with a cluster of  $N/3$  vertices and each edge with the complete bipartite graph  $K_{N/3, N/3}$ . Since  $\Gamma_3$  cannot be perfectly tiled by triangles, neither can the blown up version  $\Gamma_3(N/3)$  unless you add a single edge. Martin and Szemerédi [15] showed that Fischer's conjecture is true for  $r = 4$ . Note that in general, a tiling result for multipartite graphs does not follow from a corresponding result for arbitrary graphs. On the other hand, given a graph  $G$  of order  $nr$ , we can easily obtain (by taking a random partition) an  $r$ -partite balanced spanning subgraph  $G'$  such that  $\bar{\delta}(G') \geq \delta(G)/r - o(n)$ . Therefore a tiling result for multipartite graphs immediately gives a slightly weaker tiling result for arbitrary graphs.

The next chapter will focus on a tripartite graph and will provide a lower bound on  $\bar{\delta}(G)$ , for balanced  $G$ , in order to obtain a perfect  $K_3$ -tiling, often referring to  $K_3$  as a triangle. Here we use the absorbing lemma, though previously Magyar and Martin [14], by using Szemerédi's Regularity Lemma, were able to avoid  $\gamma$ . The advantage in using the absorbing method is that we will achieve a much smaller order graph than is necessary with the Regularity Lemma.

**Theorem 1.1.** *For any  $\gamma > 0$ , there exists  $n_0$  such that for all  $n > n_0$  the following holds: Let  $G$  be a balanced tripartite graph on  $n = 3N$  vertices with  $\bar{\delta}(G) \geq (2/3 + \gamma)N$ , then  $G$  contains a  $K_3$ -factor.*

The last chapter focuses on tiling problems in hypergraphs. We say that a hypergraph  $\mathcal{H}$  is  $k$ -uniform, also called a  $k$ -graph, if every edge in  $E(\mathcal{H})$  contains exactly  $k$  vertices. We denote the complete  $k$ -graph on  $n$  vertices by  $K_n^k$ . For a set  $T$  of size  $l < k$  in  $\mathcal{H}$ , we define  $\deg(T)$  to be the number of edges in  $\mathcal{H}$  that contain  $T$  and  $\delta_l(\mathcal{H})$  be the minimum  $l$ -degree of  $\mathcal{H}$ . For  $l = k - 1$ , we say that  $\delta_{k-1}(\mathcal{H})$  is the minimum vertex codegree of  $\mathcal{H}$ . All hypergraphs in this chapter will be 3-graphs.

**Definition 1.2.** Let  $t_l^k(n, F)$ , for all integers  $k > l \geq 1$  and  $n \in k\mathbb{Z}$ , denote the minimum  $t$  such that every  $k$ -uniform hypergraph  $\mathcal{H}$  on  $n$  vertices satisfying  $\delta_l(\mathcal{H}) \geq t$  contains a perfect  $F$ -tiling.

In their survey on the subject, Rödl and Ruciński [18] point out this result from Kühn and Osthus [10]:

$$t_2^3(n, C_4^{(3,1)}) \sim n/4,$$

where the graph  $C_4^{(3,1)}$  is the (3, 1)-cycle graph on 4 vertices.

When  $k = 2$  this is exactly the graph case and has been discussed above. For  $k \geq 3, l = k - 1$  Kühn and Osthus [11], as well as Rödl et. al. [19–21], investigated the number  $t_{k-1}^k(n, F)$ . Notably, Rödl, Ruciński and Szemerédi [20] determined  $t_{k-1}^k(n, F)$  for arbitrary  $k \geq 3$  and sufficiently large  $n$ , showing  $t_{k-1}^k(n, F) = n/2 - k + c_{k,n}$  where  $c_{k,n} \in \{3/2, 3, 5/2, 3\}$  based on the parities of  $k$  and  $n$ . Continuing this work, Pikhurko [17] provided the bounds

$$\frac{3}{4}n - 2 \leq t_l^k(n, K_4^3) \leq \frac{2 + \sqrt{10}}{6}n + O(\sqrt{n \log N}),$$

where the upper bound was also proved, independently by Keevash and Zhao (unpublished).

For the upper bound on  $t$  for  $K_k^3$ -tilings we extend an argument from Fischer [7] by introducing a weight function to handle the added complexity of the hypergraph.

**Theorem 1.3.** *Let  $\mathcal{H}$  be a 3-graph of order  $n$  with  $\delta_2(\mathcal{H}) \geq \left(1 - \frac{2}{k(k-2)}\right)n$  and  $k|n$ . Then there exists a tiling of vertex disjoint copies of  $K_k^3$  in  $\mathcal{H}$  covering all but at most  $k^2$  vertices.*

Lo and Markström [13] have a proof that extends this proof to all  $K_k^t$ -tilings, obtaining the same bound.

To show the lower bound on  $t$  we extend a construction from Pikhurko [17] to show that  $\mathcal{H}$  may not contain a  $K_k^3$ -factor.

**Proposition 1.4.** *Let  $\mathcal{H}$  be 3-graph on  $n = 2kq + r$  for integers  $k, q \geq 0$  and  $r \in \{0, k\}$ , we have*

$$\delta_2(\mathcal{H}) \geq 2(k-1)q + r - 2 \geq \left(1 - \frac{1}{k}\right)n - 2.$$

Lo and Markström [13] also extended this construction to all  $K_k^t$  and achieved an improved bound.

## Chapter 2

### PROOF OF THEOREM 1.1

Let  $\gamma > 0$  and  $n_0(\gamma)$  be the minimum positive integer satisfying the following two conditions:

$$(i) \quad 2\gamma^2 n_0^2 + \frac{5}{3}\gamma n_0^2 + 1 \geq 3\gamma n_0 + n_0$$

$$(ii) \quad 6\gamma^2 n_0^2 + 2 \geq 7\gamma n_0 + \frac{2}{3}n_0$$

Also let  $G = (V_1, V_2, V_3, E)$  be a balanced tripartite graph of order  $n = 3N$  with  $\bar{\delta} \geq (2/3 + \gamma)N$ . We prove Theorem 1.1 in three steps. First we show that for an arbitrary  $T = \{v_1, v_2, v_3\}, v_i \in V_i$ , there are many absorbing 6-sets. Next we show that  $G$  will have a near perfect tiling that misses only six vertices. Last, we will show that the final six vertices can be absorbed into the tiling.

#### 2.1 Absorbing Sets

We use Proposition 2.1 to establish an absorbing structure in  $G$  and prove that the edge density provides enough absorbing 6-sets for an arbitrary  $T$  to be added to a partial tiling. The proof follows from Lemma 10 (Absorbing Lemma) by Han et. al. [9].

**Proposition 2.1.** *For  $G$ , as in the theorem, there exists a tiling  $M$  in  $G$  of size  $|M| \leq \frac{1}{2}\gamma^2 N$  such that for every set  $W \subset V \setminus V(M)$  of size at most  $\frac{1}{2}\gamma^6 N$  there exists a tiling covering exactly the vertices in  $V(M) \cup W$ .*

*Proof.* In  $G$  we say that a set  $A = A_1 \cup A_2 \cup A_3, A_i \in \binom{V_i}{2}$ , is an absorbing 6-set for  $T$  if  $A$  spans a tiling of size 2 and  $A \cup T$  spans a tiling of size 3. Lemma 2.2 determines how many such  $A$  exist for arbitrary  $T$ .

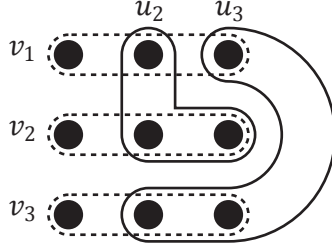


Figure 2.1. An Absorbing Structure

**Lemma 2.2.** *For every  $T$  in  $G$ , there are at least  $\frac{2}{9}\gamma^2 N^6$  absorbing 6-sets for  $T$ .*

*Proof.* Fix a set  $T$ . We wish to build the structure in Figure 2.1, so we begin by finding a triangle containing  $v_1$  but not  $v_2$  or  $v_3$ . By the degree condition,  $v_1$  has at least  $(2/3 + \gamma)N - 1$  vertices in  $V_2$  that are not  $v_2$ . Let  $u_2 \neq v_2$  be a neighbor of  $v_1$  and consider  $N_{V_3}(v_1) \cap N_{V_3}(u_2)$ . The shared neighborhood of  $v_1$  and  $u_2$  that avoids  $v_3$  must be at least

$$(2/3 + \gamma)N + (2/3 + \gamma)N - N - 1 = (1/3 + 2\gamma)N - 1$$

vertices  $u_3 \neq v_3$ . Thus, we have in total

$$((2/3 + \gamma)N - 1)((1/3 + 2\gamma)N - 1) \geq \frac{2}{9}N^2 \quad (2.1)$$

triangles that contain  $v_1$  and not  $v_2$  or  $v_3$ , as  $N \rightarrow \infty$ .

Fix one such triangle  $\{v_1, u_2, u_3\}$  and let  $U_1 = \{u_2, u_3\}$ . Now suppose we are able to choose a set  $U_2$  such that it is disjoint to  $U_1 \cup T$  and both  $U_2 \cup \{u_2\}$  and  $U_2 \cup \{v_2\}$  are triangles in  $G$ . Suppose further that we are able to choose a set  $U_3$  such that it is disjoint to  $U_1 \cup U_2 \cup T$  and both  $U_3 \cup \{u_3\}$  and  $U_3 \cup \{v_3\}$  are triangles in  $G$ . Then we call such a choice for  $U_2$  and  $U_3$  good, motivated by  $U_1 \cup U_2 \cup U_3$  being an absorbing 6-set for  $T$ , which describes the structure shown in Figure 2.1.

Focus on the number of good sets for  $U_2$ . The shared neighborhood of  $u_2$  and  $v_2$  in  $V_1$  is at least  $(1/3 + 2\gamma)N - 1$  vertices avoiding  $v_1$ . Fix a vertex  $x_1 \neq v_1$  and count how many of its neighbors in  $V_3$  are also adjacent to both  $v_2$  and  $u_2$ , while avoiding  $v_3$ . The vertices  $x_1, v_2$  and  $u_2$

will have at least  $(1/3 + 2\gamma)N + (2/3 + \gamma)N - N - 2 = 3\gamma N - 2$  common neighbors in  $V_3$  that avoid  $v_3$  and  $u_3$ . We have in all at least

$$((1/3 + 2\gamma)N - 1)(3\gamma N - 2) \geq \gamma N^2 \quad (2.2)$$

good choices for  $U_2$ . The same analysis hold for the number of choices for  $U_3$ .

Using equations (2.1) and (2.2), we see that the total number of absorbing 6-sets for  $T$  is

$$\frac{2}{9}N^2 \times (\gamma N^2)^2 = \frac{2}{9}\gamma^2 N^6.$$

□

To continue the proof of Proposition 2.1, we let  $\mathcal{L}(T)$  denote the family of all the 6-sets that can absorb the  $T$  fixed in Lemma 2.2. We know that  $|\mathcal{L}(T)| \geq \frac{2}{9}\gamma^2 N^6$ , again from Lemma 2.2. Choose a family  $\mathcal{F}$  of 6-sets by selecting each of the  $\binom{N}{2}^3$  possible 6-sets independently with probability

$$p = \frac{\gamma^3}{N^5}.$$

Then we can use the following result by Chernoff (see [1]) to determine how big  $\mathcal{F}$  is likely to be.

**Proposition 2.3.** *If  $X_i, 1 \leq i \leq n$ , be mutually independent random variables with*

$$Pr[X_i = +1] = Pr[X_i = -1] = \frac{1}{2}$$

and set

$$S_n = X_1 + \cdots + X_n.$$

Let  $a > 0$ . Then

$$Pr[S_n > a] < e^{-a^2/2n}.$$

Therefore, with probability  $1 - o(1)$ , as  $N \rightarrow \infty$  the family  $\mathcal{F}$  fulfills the following properties:

$$|\mathcal{F}| \leq 2E(|\mathcal{F}|) \leq 2\frac{\gamma^3}{N^5} \binom{N}{2}^3 \leq \frac{1}{4}\gamma^3 N \quad (2.3)$$

and

$$|\mathcal{L}(T) \cap \mathcal{F}| \geq \frac{1}{2} \mathbb{E}(|\mathcal{L}(T) \cap \mathcal{F}|) \geq \frac{1}{2} \left( \frac{\gamma^3}{N^5} \right) \times \frac{2}{9} \gamma^2 N^6 \geq \frac{1}{9} \gamma^5 N \quad (2.4)$$

Moreover we can bound the expected number of intersecting 6-sets by choosing a 6-set, a vertex in the 6-set, a second vertex in same partition and a pair of vertices from each of the other two partitions:

$$\binom{N}{2}^3 \times 6(N-1) \binom{N}{2}^2.$$

Then, the probability of choosing both sets is

$$p^2 \binom{N}{2}^3 \times 6(N-1) \binom{N}{2}^2 \leq \frac{1}{4} \gamma^6 N \quad (2.5)$$

Now, in order to upper bound the number of intersecting sets we use Markov's bound (also in [1]).

**Proposition 2.4.** *Suppose that  $Y$  is an arbitrary nonnegative random variable,  $\alpha > 0$ . Then*

$$\Pr[Y > \alpha E[Y]] < 1/\alpha.$$

Therefore, with probability at least  $1/2$

$$\mathcal{F} \text{ contains at most } \frac{1}{2} \gamma^6 N \text{ intersecting pairs.}$$

Therefore, with positive probability the family  $\mathcal{F}$  has the properties stated in (2.3), (2.4) and (2.5). Since some of the 6-sets will not absorb any  $T$  and some will intersect each other, we delete all of these undesired 6-sets in the family  $\mathcal{F}$  to get a subfamily  $\mathcal{F}'$  consisting of pairwise disjoint absorbing 6-sets which satisfies

$$|\mathcal{L}(T) \cap \mathcal{F}'| \geq \frac{1}{9} \gamma^5 N - \frac{1}{2} \gamma^6 N \geq \frac{1}{2} \gamma^6 N.$$

Finally, the thinned out family  $\mathcal{F}'$  consists of pairwise disjoint absorbing 6-sets and  $G[V(\mathcal{F}')] contains a perfect tiling  $M$  of size at most  $\frac{1}{2} \gamma^3 N$ . Also, for any subset  $W \subset V \setminus V(M)$  of size  $\frac{1}{2} \gamma^6 N$$



we can partition  $W$  into sets of size 3 and successively absorb them using a different absorbing 6-set each time. This gives us a tiling that covers exactly the vertices in  $V(\mathcal{F}') \cup W$ .  $\square$

## 2.2 Complete Tiling

To complete the proof of the theorem, we find in  $G$  an absorbing family  $M$  guaranteed by Proposition 2.1. We let  $G' = G - V(M)$  and observe that

$$\bar{\delta}(G') \geq (2/3 + \gamma)N - \frac{3}{2}\gamma^3 N \geq \frac{2}{3}N \geq \frac{2}{3}N'$$

where  $N'$  is the number of vertices in each partition set of  $G'$ . Notice further that  $G'$  is still balanced and we can apply Proposition 3.2 in Fischer [7] to find an incomplete tiling in  $G'$ .

**Proposition 2.5.** *If  $G$  is a tripartite graph with vertex partitions  $V_1, V_2$  and  $V_3$  of size  $N$ , such that each vertex in any partition has at least  $\frac{2}{3}N$  neighbors in each of the other partitions, then  $G$  contains  $N - 2$  disjoint triangles.*

This proposition gives us an almost perfect tiling of  $G'$ , leaving only a set  $W$  containing 6 vertices uncovered. By Proposition 2.1 we can divide  $W$  into sets of 3 and use  $M$  to absorb each triple and complete the perfect tiling on  $G$ .

## Chapter 3

### PROOFS ON 3-GRAPHS

In this chapter we provide a minimum degree condition that guarantees an almost perfect tiling of a 3-graph  $\mathcal{H}$  that misses at most  $k^2$  vertices. Next we will provide a construction that shows that if the minimum degree condition is too small, we cannot guarantee a perfect tiling of  $\mathcal{H}$ .

#### 3.1 Proof of Theorem 1.3

This proof is adapted from the proof of Lemma 6.1 by Pikhurko [17] which adapts the proof of Theorem 2.1 by Fischer [7].

*Proof.* Let  $\mathcal{H}$  be a 3-graph on  $n$  vertices with  $\delta_2(\mathcal{H}) \geq \left(1 - \frac{2}{k(k-2)}\right)n$  and  $k|n$ . Begin with a partition  $\mathcal{P}$  of the vertex set  $V(\mathcal{H})$  into sets of size  $k$ ,  $V_1, \dots, V_{n/k}$ . Let  $G_i$  be the largest complete graph in  $V_i$ . If  $V_i$  is an independent set, we define  $|G_i| = 2$ . Denote by  $w : \{2, \dots, k\} \rightarrow \mathbb{R}$  the function defined by  $w(2) = 0$  and  $w(j+1) - w(j) = 1 - \frac{1}{k^j}$  for  $2 \leq j \leq k-1$ . We say that  $w(\mathcal{P})$ , the weighting of  $\mathcal{P}$ , is  $\sum_{1 \leq j \leq n/k} w(|G_j|)$ . Assume that  $\mathcal{P}$  is chosen such that  $w(\mathcal{P})$  is maximal. We will now show that for each weight class  $2 \leq i \leq k-1$  there are at most  $k-1$  sets  $V_j$  in  $\mathcal{P}$  with  $|G_j| = i$ . Suppose, for a contradiction, that  $|G_1| = \dots = |G_k| = i < k$ . Since  $|G_j| < k$  for  $1 \leq j \leq k$  we can find at least one  $v_j \in V_j \setminus G_j$ . Now, for  $1 \leq j \leq k$  and vertex  $v \notin V_j$ , we say the pair  $(v, j)$  is a connection if and only if  $\{v\} \cup G_j$  spans a complete hypergraph. If there are any connections  $(v, j)$  with  $v \in V_1 \cup \dots \cup V_k$  then switching  $v$  with any vertex  $v_j$  will result in a new partition  $\mathcal{P}'$ . Note that since

$$1 - \frac{1}{k^i} \geq 1 - \frac{1}{k^{i-1}}$$

we have

$$w(i+1) - w(i) \geq w(i) - w(i-1)$$

which is

$$w(i+1) + w(i-1) \geq 2w(i)$$

and we immediately provide a contradiction to  $w(\mathcal{P})$  being maximal. Thus, we can assume there are no connections with  $v \in V_1 \cup \dots \cup V_k$  and  $1 \leq j \leq k$ .

Using the condition on  $\delta_2(\mathcal{H})$ , for  $1 \leq j \leq k$  we can determine a lower bound on the number of connections there are by double counting the number of adjacencies among the  $G_j$ 's. An arbitrary pair of vertices in  $G_j$  is adjacent to at least  $\delta_2(\mathcal{H})$  vertices. If we let  $c$  be the number of connections to  $G_j$  then

$$\binom{i}{2} \delta_2(\mathcal{H}) \leq \binom{i}{2} c + \left( \binom{i}{2} - 1 \right) (n - c)$$

and

$$c \geq \binom{i}{2} \delta_2(\mathcal{H}) - \left( \binom{i}{2} - 1 \right) n \geq \frac{(k-i)n}{k}$$

where the last inequality is true since  $i < k$ .

Now there are at least  $(k-i)n$  connections  $(v, j)$  with  $v \notin V_1 \cup \dots \cup V_k$  and  $1 \leq j \leq k$ . Since  $n > k$  we can choose  $V'_j$  such that there are more than  $k(k-i)$  connections  $(v', j)$  for  $v' \in V'_j$  and  $1 \leq j \leq k$ . Consider the bipartite graph  $B$  with parts  $\{G_1, \dots, G_k\}$  and  $V'_j$  whose edge set consists of those pairs that make a connection. Since  $B$  has at least  $k(k-i)$  edges, the König-Egerváry Theorem (see [4] Theorem 8.32) shows that  $B$  contains a matching of size at least  $k-i+1$ . Now by moving  $v'_j$  to  $V_j$  for  $1 \leq j \leq k-i+1$  and  $\{v_1, \dots, v_{k-i+1}\}$  to  $V'_j$ , see Figure 3.1,  $w(\mathcal{P})$  increases by

$$\begin{aligned} & (k-i+1)(w(i+1) - w(i)) - (w(|G'_j|) - w(\max\{2, |G'_j| - k + 1 + i\})) \\ & \geq (k-i+1) \left( 1 - \frac{1}{k^i} \right) - \left( k+1-i - \frac{k-i+1}{k} \right) \\ & = \frac{(k^i - 1)(k-i+1)}{k^{i+1}} > 0 \end{aligned}$$

a contradiction.

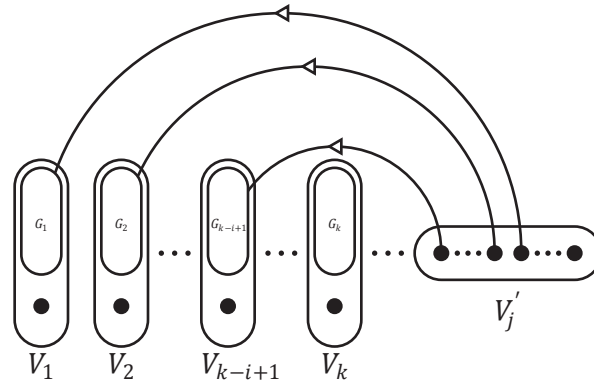


Figure 3.1. Vertices making a connection from  $V'_j$

□

### 3.2 Proof of Proposition 1.4

We now provide a construction that proves that the codegree of  $\mathcal{H}$  must be larger than  $(1 - 1/k)n - 2$  if we are to be guaranteed a perfect tiling.

*Proof.* For  $n = 2kq + r$ , if  $r = k$  let  $a_0 = 2q + 1$ . Otherwise we let  $a_0$  be either  $2q + 1$  or  $2q - 1$ , with both choices giving the same bound. Partition  $V(\mathcal{H}) = A_0 \cup A_1 \cup \dots \cup A_{k-1}$  into parts of sizes  $a_0 + a_1 + \dots + a_{k-1} = n$ , where  $a_1, \dots, a_{k-1}$  are nearly equal, that is  $|a_i - a_j| \leq 1$  for  $1 \leq i < j \leq k - 1$ . Let  $\mathcal{H}$  be the 3-graph on  $n$  vertices whose edge set consists of all triple excluding any that satisfy one of the following (mutually exclusive) properties:

- (i) have exactly three vertices in  $A_0$
- (ii) have one vertex in  $A_0$  and two vertices in  $A_i$  for some  $1 \leq i \leq k - 1$
- (iii) intersect each of  $A_1, A_2$  and  $A_3$ .

Figure 3.2 shows examples of edges that are excluded from  $\mathcal{H}$ . To see why there can be no  $K_k^3$ -tiling, consider any  $K_k^3$ -subgraph  $K$  of  $\mathcal{H}$ . By Property (i),  $K$  cannot intersect  $A_0$  in more than two vertices. Suppose that  $K$  intersects  $A_0$  in exactly one vertex and avoids at least one partition. Then by the pigeon hole principle there is a partition  $A_i$  for  $1 \leq i \leq k - 1$  that contains at least

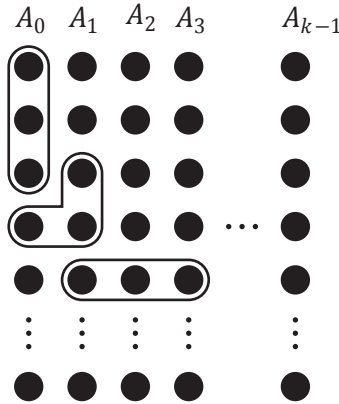


Figure 3.2. Examples of Edges Not Allowed

two vertices of  $K$ . Property (ii) forbids the edge spanning the vertex in  $A_0$  along with any pair in  $A_i$ . So if  $K$  is to intersect  $A_0$  in exactly one vertex,  $K$  must also intersect every other partition in exactly one vertex. By property (iii), the edge with a vertex in  $A_1, A_2$  and  $A_3$  is forbidden, so  $K$  cannot intersect  $A_0$  in one vertex in this manner either.

Therefore every  $K_k^3$ -subgraph of  $\mathcal{H}$  has an even number of vertices in  $A_0$ . This makes a perfect tiling impossible, since  $|A_0| = 2q \pm 1$ , which is odd.

A case by case analysis gives the desired bound.

**Case 1** Two vertices in  $A_0$  are in an edge with every vertex in  $A_i$  for  $1 \leq i \leq k-1$ , so the codegree is  $\frac{k-1}{k}n$ ;

**Case 2** One vertex in  $A_0$  and one vertex in  $A_i$  for  $1 \leq i \leq k-1$  are in an edge with every other vertex in  $A_0$  and every vertex in  $A_j$  for  $j \neq i$  and  $1 \leq j \leq k-1$ , so the codegree is  $\frac{k-1}{k}n - 1$ ;

**Case 3** Two vertices in  $A_i$  for  $1 \leq i \leq k-1$  are in an edge with every other vertex in  $A_i$  and every vertex in  $A_j$  for  $j \neq i$  and  $1 \leq j \leq k-1$ , so the codegree is  $\frac{k-1}{k}n - 2$ ;

**Case 4** One vertex in  $A_i$  and one vertex in  $A_j$  for  $i, j \in [3]$  and  $i \neq j$  are in an edge with every vertex in  $A_0$ , every other vertex in  $A_i$  and  $A_j$  and every vertex in  $A_\ell$  for  $4 \leq \ell \leq k-1$ , so the codegree is  $\frac{k-1}{k}n - 2$ ;

**Case 5** One vertex in  $A_i$  for  $i \in [3]$  and one vertex in  $A_j$  for  $4 \leq j \leq k - 1$  are in an edge with every other vertex of  $\mathcal{H}$ , so the codegree is  $n - 2$ ;

**Case 6** Two vertices in  $A_i$  for  $4 \leq i \leq k - 1$  are in an edge with every other vertex of  $\mathcal{H}$ , so the codegree is  $n - 2$ .

We take the minimum of these codegrees, which is  $\frac{k-1}{k}n - 2$ .

□

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