

DISCLAIMER

This report was prepared as an account of work sponsored by an agency of the United States Government. Neither the United States Government nor any agency thereof, nor any of their employees, make any warranty, express or implied, or assumes any legal liability or responsibility for the accuracy, completeness, or usefulness of any information, apparatus, product, or process disclosed, or represents that its use would not infringe privately owned rights. Reference herein to any specific commercial product, process, or service by trade name, trademark, manufacturer, or otherwise does not necessarily constitute or imply its endorsement, recommendation, or favoring by the United States Government or any agency thereof. The views and opinions of authors expressed herein do not necessarily state or reflect those of the United States Government or any agency thereof.

DISCLAIMER

Portions of this document may be illegible in electronic image products. Images are produced from the best available original document.

4	Table of	A	General Information	
\bigcirc	Table of Contents	ß	Biographies & Abstracts	
-	Keynote Tutorial	Ç	Dave Parnas: Design Through Documentation	••
		D	W1: Natural Language Modeling	
		E	X1: Defining Software Processes	
	Tutorials		Y1: How the NWC Handles Software as Product	
	April 1, 1997	G	W2: Writing Testable Software Requirements	
		Н	X2: Using COTS Software in Development Projects	
			Y2: Software Inspection Process Overview	
-	Keynote Address	J	Capers Jones: Software Quality for 1997	
		К	AI: Software Management	
		L	B1: Software Testing	
\bigcirc	General Sessions	M	C1: Software Quality for Scientific Applications	
Ċ	April 2, 1997		A2: Software Engineering Processes	
	_		B2: Internet WEB Applications	
		P	A3: Software Process Improvement I	
		Q	B3: High Integrity / Formal Methods I	
-	General Sessions April 3, 1997	R	A4: Software Process Improvement II	
		S	B4; High Integrity / Formal Methods II	
		Т	A5: Software Quality: Experiences & Year 2000	
		U	B5: Software Standards for Quality Engineering	
	Closing Session	V	Wrapup and Awards	
		W	Pre-Forum Registrations	
		X	Final Forum Attendance List	
Č.		Y	Notes	1015 1015 1015 1015 1015 1015 1015 1015
		Z	Notes	
		&anera.	Ready Index® Indexing System	



General Information

- Summary Schedule
- Forum Committee & Program Committee
- History of the Software Quality Forum
- Software Quality Assurance Subcommittee
- Forum Awards, Proceedings, and Participating Organizations
- Tours of National Atomic Museum and RMSEL
- No Host Dinner El Pinto Restaurant
- Bus Schedule for Social, Tours, and No-Host Dinner
- Maps: Local Area and Forum Site

SOFTWARE QUALITY FORUM April 1-3, 1997 CONFERENCE SUMMARY

11:00 + 01:00 pm 11:00 + 01:00 pm 22: high 10:00 - 03:00 pm 03:00 + 03:15 pm 03:00 + 03:15 pm 03:20 - 05:30 pm 03:20 - 05:30 pm 03:30 - 05:30 pm 03:30 - 05:00 pm 10:15 - 11:45 pm 10:15 - 11:45 pm 11:45 - 01:20 pm 11:45 - 01:20 pm 01:20 - 03:00 pm 01:20 - 03:00 pm 03:00 + 03:15 pm	Keynote Tutorial TIC Auditorium Truck Z - Keynote Tutorial SNL Bilg \$23, Brazeway icuos of Cruzel Sethers of Press, Millister University Welcoming Remarks TIC Auditorium Keynote Address TIC Auditorium Keynote Address TIC Auditorium A TYC Auditorium Al; Seterer Mangemet Car: Den Schling, ASTIMET	REGISTRATION/CONTINENTAL BI Design Through Documentation: The Pith is Software Quality Dr. David Farmer, McMaster University LUNCH - ON YOUR OWN Track W SNL Bidg 222, Room A WI: Heard Language labeling Re, Jain Shap, SH, BREAK - TTC LOBBY WE Wrong Temble SW Requirements Dr. Dasgae Reak, 594 Social Hour and Bird's of a Feather National Adomic Musicities National Adomic Musicities Registration, April 2, 1997 REGISTRATION/CONTINENTAL Bit Make Blackladge, Forem Chair Far Whiteward, SNL Director John Crewford, SNL Director John Crewford, SNL Director John Crewford, SNL Director John Crewford, SNL Director BREAK - TTC LOBBY Track B SNL Bidg 322 Roome A&B BREAK - TTC LOBBY LUNCH - ON YOUR OWN	Track X SNL Bidg \$22. Room B 30: Defining Selferen Process Sr. Genal McDaniki, Stil, Cantradar 22. Ung COTO Astronom in Development Projects Is Col Manay Crowley, USAP Philips Laboratory Meet in TSC Lobby - Round Trip Tran	Barry Late and Rundy Dollar, 24,					
01:00 - 03:00 pm 03:00 - 03:15 pm 03:00 - 03:15 pm 03:15 - 05:15 pm 05:20 - 06:30 pm 05:20 - 06:30 pm 05:30 - 06:20 pm 09:00 - 10:00 pm 10:15 - 11:45 pm 10:15 - 11:45 pm 11:45 - 01:20 pm 01:20 - 03:00 pm 01:20 - 03:00 pm 03:00 + 03:15 pm 03:00 + 03:15 pm	TTC Auditorium Truck Z - Keynote Totorial SNL Brig \$23, Branzowsy cours of Crown Schwarswy Rail and Decorption McManare University Welcoming Remarks TTC Auditorium Keynote Address TTC Auditorium Keynote Address TTC Auditorium All: Schwar Mangemei Cher Den Schwar, ASTMAT	The Path to Software Quality Dr. David Farmer, McMinister University LUNCH - ON YOUR OWN Track W SNL Bidg 222, Room A WI: Humai Lagrage Modeling So. Lein Sharp, Sel. BREAK - TTC LOBBY WP Wrange Tambie SW Requirements Cr. Danges Rook, SW Social Hour and Births of a Feather National Alonne: Museeurs National Alonne: Museeurs National Alonne: Museeurs Weddrefficially, April 2, 1997 REGISTRATION/CONTINENTAL BI Make Blackladge, Forum Chair Eart Whitenson, DOE/AL Director John Crowford, SNL Elecanive VP Software Quality for 1997 - What Works and What Doesn't? Cancers Jones, Chairman Software BIREAK - TTC LOBBY Track B SNL Bidg \$22 Rooms A&B BIREAK - TTC LOBBY	SNL Bidg 172, Room B 21: Defaug Eddress Process 22: Dang COTE Johann & Development Projects 12: Conference of Neurophysics 12: Conference Room C Track C TPC Conference Room C C1: SW Queby for Security Againston	TTC Conference Room T1::Bar the Nove Hardles Software in Free Dand Views, Pesce					
01;00 - 03:00 pm 21: basis 03:00 - 03:15 pm 21: State 03:15 - 05:15 pm 21: State 03:20 - 06:30 pm 21: State 05:20 - 06:30 pm 21: State 07:30 - 08:20 am 21: State 09:00 - 10:00 am 21: State 10:00 - 10:15 am 21: State 10:15 - 11:45 am 21: State 11:45 - 01:20 pm 21: State 01:20 - 03:00 pm 21: State 03:00 - 03:00 pm 21: State 03:00 - 03:00 pm 21: State 03:00 - 03:00 pm 21: State	SNL Bidg \$23, Brazzowry cours of Croupl Sectorem of Parene, McMasters University Radi and Decorpton Ad Parene, McMasters University Wielcouring Remarks TTC Authoritan Keynotic Address TTC Authoritan Keynotic Address TTC Authoritan Al; Sectorer Mangement Char: Den Schling, ASTMACT	Track W SNL Bidg 222, Room A WI: Hannel Langung Modeling De, Min Sharp, SH. BREAK - TTC LOBBY WI: George Temble SW Requirements Cr. Dasgee Kark SH. Social Hour and Births of a Feather National Alornic Musleum National Alornic Musleum Wednesselay, April 2, 1997 REGISTRATION/CONTINENTAL BI Malo Blackladge, Forem Chair Earl Whiteness, DOE/AL Director John Crewford SML Becaure VP Software Quality for 1997 - What Works and What Doem?* Cancers Jones, Chairman Software BREAK - TTC LOBBY Track B SNL Bidg \$22 Roomes A&B BREAK - TTC LOBBY	SNL Bidg 172, Room B XI: Defang Sidowa Promos 2. Genil McDanki, S4, Castraire XI: Useg COTS Johann & Drokenment Projecto II: Col Nany Crowley, USAP Ridge Laboratory Meet in TTC Lobby - Round Top Tran XEAK FAST - TTC LOBBY Track C Track C TTC Conference Room C CI SW Quity for Security Againston	TTC Conference Room T1::Bar the Nove Hardles Software in Free Dand Views, Pesce					
01;00 - 03:00 pm 21: basic 03:00 - 03:15 pm 0: Dec Device 03:15 - 05:15 pm 21: State 03:20 - 06:30 pm 21: State 05:20 - 06:30 pm 21: State 09:00 - 10:00 ant 10:00 + 10:15 am 10:00 + 10:15 am 10:15 - 11:45 am 10:15 - 11:45 am 11:45 - 01:20 pm 01:20 - 03:00 pm 10:00 + 10:15 am	SNL Bidg \$23, Brazzowry cours of Croupl Sectorem of Parene, McMasters University Radi and Decorpton Ad Parene, McMasters University Wielcouring Remarks TTC Authoritan Keynotic Address TTC Authoritan Keynotic Address TTC Authoritan Al; Sectorer Mangement Char: Den Schling, ASTMACT	SNL Bidg 222, Room A WI: Named Langung Matching (a. Main Sharp, Sel. BREAK - TTC LOBBY WI: Wange Trankle SW Requestories Cr. Darger Reach GRU Social Hour and Bards of a Feather Nahoral Alorme Museum Nahoral Alorme Museum WEDDETGIAY, April 2, 1997 REGISTRATION/CONTINENTAL BI Make Blackladge, Forem Chair Eart Whiteness, DOE/AL Director John Crowford, SNL Excame VP Software Quality for 1997 - What Works and What Doesn's Cancers Jones, Chairman Software BREAK - TTC LOBBY Track B SNL Bidg \$22 Rooms A&B B: Software Turny, Cancer Lany Roda, Futur	SNL Bidg 172, Room B XI: Defang Sidowa Promos 2. Genil McDanki, S4, Castraire XI: Useg COTS Johann & Drokenment Projecto II: Col Nany Crowley, USAP Ridge Laboratory Meet in TTC Lobby - Round Top Tran XEAK FAST - TTC LOBBY Track C Track C TTC Conference Room C CI SW Quity for Security Againston	TTC Conference Room T1::Bar the Nove Hardles Software in Free Dand Views, Pesce					
01:00 - 03:00 pm Dt Die 03:00 - 03:15 pm Dt Die 03:15 - 05:15 pm Dt Die 03:15 - 05:15 pm Dt Die 05:20 - 06:30 pm Dt Die 05:20 - 06:30 pm Dt Die 05:20 - 06:30 pm Dt Die 05:30 - 06:00 pm Dt Die 00:00 - 10:00 pm Dt Die 10:15 - 11:45 sm Dt Die 10:15 - 11:45 sm Dt Die 01:20 - 03:00 pm Dt D	Al Parent, Mildester Charactery Al Parent, Mildester Charactery Welcouring Remarks TTC Auditorium Keynote Address TTC Auditorium Track A TYC Auditorium Al; Science Mangamed Char Den Schling, ASTMACT	 (iv, sein Sharp, Sei, IRESAK - TTC LOBBY WP Wrong Tumble SW Requirements Despectively State SW. Social Hour and Bird's of a Feather National Atomic Musicities Weidnersday, April 2, 1997 REGISTRATION/CONTINENTAL BI Make Bleckladge, Forum Chair Fart Wintensen, DOE/AL Director John Crewford, SNL Excounte VP Software Quality for 1997 - What Works and What Doem?* BREAK - TTC LOBBY Track B SNL Bidg \$22 Rooms A&B B:: Softwar Tump, Court (my Rode, Future 	Centre in TTC Lobby - Round Trajector ColVery Coviey, USAP Printipe Laboratory Meet in TTC Lobby - Round Trajectory Centre C TTC Conference Rotes C CI SW Quity for inset for Agriculture	Dend Views, Pescs					
03:15 - 05:15 pm 0. Den	Millionning Remarks Wielcoming Remarks TTC Auditorium Keynote Address TTC Auditorium Track A T?C Auditorium Al; Schwer Mangemei Cher: Den Schwer, ASTHET	WE Wrang Temble SW Requesters Dr. Dasget Kork, SW Social Hour and Birtls of a Feather National Alonne Musicura WEODERICAY, April 2, 1997 REGISTRATION/CONTINENTAL BI Mules Bleckladge, Forem Chair Earl Whitensen, DCE/AL Director John Crewford, SML Excounte VP Software Quality for 1997 - What Works and What Doem?? Cancer Jones, Chairman Software BREAK - TTC LOBSY Track B SNL Bidg \$22 Rooms A&B B:: Software Temp Conce Lary Rode, Patter	Li Col Many Crowley, USAP Ridge Laboratory Meet in TTC Lobby - Round Trip Tran EAKPAST - TTC LOBBY TEAKPAST - TTC LOBBY	Track D					
03:15 - 05:35 pm Cr. Devi 05:20 - 06:30 pm Cr. Devi 77:30 - 08:20 em Cr.	Millionning Remarks Wielcoming Remarks TTC Auditorium Keynote Address TTC Auditorium Track A T?C Auditorium Al; Schwer Mangemei Cher: Den Schwer, ASTHET	Cr. Dagae Kark 594. Social Hour and Births of a Feather National Alonne: Museum Wednesday, April 2, 1997 REGISTRATION/CONTINENTAL BI Males Blackledge, Forem Chair Bart Whitenen, DOE/AL Director John Crewford, SML Brecaure VP Software Quality for 1997 - What Works and What Doem?* Cancers Jones, Chairman Software BREAK - TTC LOBBY BREAK - TTC LOBBY BREAK - TTC LOBBY BREAK - TTC LOBBY	Li Col Many Crowley, USAP Ridge Laboratory Meet in TTC Lobby - Round Trip Tran EAKPAST - TTC LOBBY TEAKPAST - TTC LOBBY	Track D					
67.30 - 08:30 am 08:30 - 69.00 am 09.00 - 10.00 am 10.00 - 10:15 am 10:15 - 11:45 am 11:45 - 01:30 pm 01:30 - 03:00 pm 01:30 - 03:00 pm 03.00 - 63:15 pm 03.00 - 63:15 pm	Welcoming Remarks TTC Auditorium Knywite Address TTC Auditorium TC Auditorium Track A TPC Auditorium Al: Setwart Mangement Care Den Schling, ASTIMET	National Alonne Musicina Wednesday, April 2, 1997 REGISTRATION/CONTINENTAL BI Mulce Blockladge, Forem Chair Eart Whitenen, DCE/AL Director John Crowford, SNL Excentre VP Software Quality for 1997 - What Works and What Doem?* Cancer Jones, Chairman Software BREAK - TTC LOBSY Track B SNL Bidg \$22 Rooms A&B Bit Software Turing Conce Lary Rode, Panton	EAKPAST - TTC LOBBY Track C TTC Conference Room C Cl SW Quity for Security Agriculture	Track D					
57.30 - 08:30 am 08:30 - 69.00 am 10:00 - 10:00 am 10:00 - 10:15 am 10:15 - 11:45 am 11:45 - 01:30 pm 01:30 - 03:00 pm 01:30 - 03:00 pm 03:00 - 63:15 pm 03:00 - 63:15 pm	Welcoming Remarks TTC Auditorium Knywite Address TTC Auditorium TC Auditorium Track A TPC Auditorium Al: Setwart Mangement Care Den Schling, ASTIMET	REGISTRATION/CONTINENTAL BI Make Blackledge, Forem Chair Earl Whiterson, DOE/AL Director John Crewford, SML Breenawe VP Software Quality for 1997 - What Works and What Doem?? Canons Jones, Chairman Software BREAK - TPC LOBBY Track B SNL Bidg \$22 Rooms A&B B: Software Turing Court (any Rode, Paster	EAKPAST - TTC LOBBY Track C TTC Conference Room C C: SW Quity for Security Againston	Track D					
02:30 + 69.00 mm ; 09.00 - 10.00 mm ; 10.00 + 10:15 mm 10:15 - 11:45 mm 11:45 + 01:20 pm 01:30 - 03:00 pm 01:30 - 03:15 pm 03.00 + 03:15 pm (3):15 - 04:45 pm	TTC Auditorium Keynote Address TTC Auditorium TTC Auditorium Track A TYC Auditorium Al: Sciwer Mangemet Carr Den Schlier, ASTMAT	Malco Blockladge, Forem Chair Earl Whitenen, DOE/AL Director John Crowford, SNL Biesnawe VP Software Quality for 1997 - What Works and What Doem?? Carons Jones, Chairman Software BREAK - TPC LOBBY Track B SNL Bidg \$22 Rooms A&B Bit Software Tuting, Char; Lary Roda, Paster	Track C TTC Conference Room C C: SW Quity to Security Applications						
09,00 - 10.00 ant 10.00 - 10:15 mm 10:15 - 11:45 mm 11:45 - 01:20 pm 01:30 - 03:00 pm 01:30 - 03:00 pm 03.00 - 03:15 pm 03.00 - 03:00 pm	TTC Auditorium Keynote Address TTC Auditorium TTC Auditorium Track A TYC Auditorium Al: Sciwer Mangemet Carr Den Schlier, ASTMAT	Earl Whitewan, DOE/AL Director John Crewford, SNL Breasure VP Software Quality for 1997 - What Works and What Doem?* Canons Jones, Chairman Software BREAK - TTC LOBBY Track B SNL Bidg \$22 Rooms A&B B: Software Turing Char; (any Rode, Paster	TTC Conference Room C						
10.00 + 10:15 mm 10:15 - 11:45 mm 11:45 + 01:20 pm 01:30 - 03:00 pm 03:00 + 03:15 pm 03:00 + 03:15 pm 03:15 - 04:45 pm	TTC Audizonium Track A TTC Audizonium Al; Sciwer: Management Case: Den Scheing, ASTFMRT	end Whet Doem't? Canons Jones, Chaimen Software BREAK - TTC LOBBY Track B SNL Bidg \$22. Rooms A&B Bit Softwar Taing Char; Lary Roda, Paster	TTC Conference Room C						
10:15 - 11:45 sam 11:45 - 01:30 pm 01:30 - 03:00 pm 01:30 - 03:00 pm 03:00 - 03:15 pm 03:15 - 04:45 pm 03:15 - 04:45 pm	T?C Auditorium Al: Solver Mangemet Char: Den Schlier, ASTMAT	BREAK - TTC LOBBY Track B SNL Bidg \$22 Rooms A&B Bit Software Testing Char; Larry Rode, Pester	TTC Conference Room C						
11:45 - 01:30 pm 01:30 - 03:00 pm 03:00 - 03:15 pm 03:15 - 04:45 pm 03:15 - 04:45 pm	T?C Auditorium Al: Solver Mangemet Char: Den Schlier, ASTMAT	SNL Bidg \$22 Room: A&B Bit Software Tuting Court (any Roda, Paster	TTC Conference Room C						
11:45 - 01:30 pm 01:30 - 03:00 pm 03:00 - 03:15 pm 03:15 - 04:45 pm 03:15 - 04:45 pm	Cher, Dox 5 dollary, A STIMAT	Çanı; Çary Roân, Pains		Benit of a Feature / Networking					
01:30-03:00 pm 03:00-03:15 pm 03:15-04:45 pm 03:15-04:45 pm	Ind	LUNCH - ON YOUR OWN							
01:50-03:00 pm 03,00-03:05 pm 03:15 - 04:45 pm 03:15 - 04:45 pm 02:00 - 07:00 pm	1st A	LUNCH - ON YOUR OWN							
01:50-03:00 pm 03,00-03:05 pm 03:15 - 04:45 pm 03:15 - 04:45 pm 02:00 - 07:00 pm	TTC Autoorium	Track B SNL Bidg \$22, Roome A&B	Track C Tours	Insth D Tours					
03-1.5 - 04:45 pen	Al. Salaran Espectrag Prosess Crass, Kuthices Carol, DOS54Q	EX: Internal WEB Applications Claury Rays Brown, LMES 0402,	kolanis Liik Report in TTC Litky Maa in TTC Litky kaan 134 pa	Hannes Access Masters Register as TTC Lobby Meet as TTC Lobby before 1:38 per					
Cc.00 + 07.00 pm		BREAK + TTC LOBBY							
•	A3 Solven (room legerment) E Char (dire (solver, ASTMAT		El Migh kompriy / Formal Marbock J Chair Delic Pancy, SNL Men et TTC Lebby Sefert J 15 pr.						
		Thursday, April 3, 1997.		in the set of the					
08,00 - 08,30 em		LOBBY							
	Truck A Truck B TTC Auditorium SNL Bidg 872, Rooms A&B		Tack C	Track D					
08:30 - 10.00 am	AA. Saltaun Jason Inpervisian I Caus, Jula Han, AWS UK	Bit: Kigh Istagray / Fatural Methods B Chart: Larry Dalan, 594.	Darks of a Franker / Hernordung.	Berbol's Forder / New Yorking					
10.00- 10.15 am	BREAK + TTC LOBBY								
10:35 - 13:45 am	Saltum Qalay: Equators & Y2C Class Calay Kala, ASTMET	ES: 517 Sundards for Quality Engineering Chart Proty Techne, SNC	Inclusion of a Practice / Instructions	Bade of a Faster / Materiatury					

_;

_

.

.

7

.

٧.,

Forum Committee

General Chair:

Mike Blackledge, Sandia National Laboratories, mablack@sandia.gov

Tutorials and Workshops:

Dave Peercy, Sandia National Laboratories, depeerc@sandia.gov

Planning Committee:

Arrangements:

Lorraine Baca, Sandia National Laboratories Ray Berg, Sandia National Laboratories Dwayne Knirk, Sandia National Laboratories Patty Trellue, Sandia National Laboratories Gary Echert, DOE - Albuquerque Office

Theresa Griego, Sandia National Laboratories

Program Committee

Mike Blackledge, Sandia National Laboratories Patty Trellue, Sandia National Laboratories

Faye Brown, Martin Marietta Energy Systems, Y-12 Plant Kathleen Canal, DOE Headquarters John Cerutti, Los Alamos National Laboratory Orval Hart, Los Alamos National Laboratory Mike Lackner, AlliedSignal Federal Manufacturing and Technologies, Kansas City Plant Dave Peercy, Sandia National Laboratories Larry Rodin, Mason & Hanger, Pantex Plant Don Schilling, AlliedSignal Federal Manufacturing and Technologies, Kansas City Plant Pat Tempel, Sandia National Laboratories David Vinson, Mason & Hanger, Pantex Plant

History of the Software Quality Forum

The Software Quality (SQ) Forum was established by the Software Quality Assurance Subcommittee as an opportunity for all those involved in implementing SQA programs to meet and share ideas and concerns. The SQ Forum is open to the public. Participation from managers, quality engineers, and software professionals provides an ideal environment for identifying and discussing the many issues and concerns raised by the Forum attendees and speakers. The interaction provided by the Forum contributes to the realization of a shared goal - high quality software product.

Topics presented at the SQ Forum generally include: testing, software measurement, software surety, software reliability, SQA practices, assessments, software process improvement, certification and licensing of software professionals, CASE tools, software project management, inspections, and management's role in ensuring SQA.

The Software Quality Forum is held every three years; past Forums are identified below.

Date	Site
Spring 1988	Sandia National Laboratories
Spring 1991	AlliedSignal Aerospace Kansas City Division
Spring 1994	Lawrence Livermore National Laboratory

Software Quality Assurance Subcommittee

The Software Quality Assurance Subcommittee (SQAS) serves as a Technical Advisory Group on software engineering and quality initiatives and issues for the Department of Energy's Quality Managers. The Quality Manager at each DOE site has the opportunity to select one Primary and one Alternate representative to the SQAS.

The Subcommittee grew out of a Software Quality Assurance Information Exchange Forum which was held in March of 1988 at Sandia National Laboratories. The Subcommittee provides a continuing forum for the exchange of information and work issues in the area of software quality engineering.

For additional information about the SQAS, visit our web site at:

http://www.pantex.com/sqas/sqas.htm

Forum Awards

The Forum Program Committee would like to recognize those presenters who, through their tutorial or technical presentation, have made a significant contribution to the success of the Forum. A Best Tutorial and Best Presentation award will be presented at the Forum Wrap-up session on Thursday, April 3. Selection of recipients for the Awards will be determined in two parts:

- technical content, scored by the Forum Committee
- delivery and usefulness, scored by attendees

Forum Proceedings

Forum Proceedings will include abstracts and presentation materials for all technical presentations, presenter biographies, tutorial materials, and final Forum program information. Forum Proceedings will be distributed at the Forum with the registration packets. Additional Forum Proceedings can be purchased at the Registration Desk in the TTC Lobby,

Participating Organizations

AlliedSignal, Federal Manufacturing and Technologies, Kansas City Plant (AS/FM&T) Atomic Weapons Establishment, United Kingdom (AWE UK) Department of Energy, Albuquerque Office (DOE/AL) Department of Energy, Headquarters (DOE/HQ) Lawrence Livermore National Laboratory (LLNL) Locleneed Martin Energy Systems, Oak Ridge, Y-12 Plant (LMES/OR) Los Alamos National Laboratory (LANL) Mason & Hanger, Pantex Plant (Pantex) McMaster University, Communications Research Laboratory, Canada (MU/CRL) New Mexico State University (NMSU) Sandia National Laboratories (SNL) Software Productivity Research (SPR) United States Air Force, Phillips Laboratory (USAF/Phillips) Westinghouse, Savannah River Site (SRS) Pioneer Technologies (Pioneer)

National Atomic Museum Tour

Operated by the Department of Energy, The National Atomic Museum contains a large collection of declassified nuclear technology. Since its opening in 1969, the objective of the National Atomic Museum has been to provide a readily assessable repository of educational materials, and information on the Atomic Age.

Prominently featured in the museum's high bay is the story of the Manhattan Engineer District, the unprecedented 2.2 billion dollar scientific-engineering project that was centered in New Mexico during World War II.

A portion of the Museum is devoted to exhibits on the research, development, and use of various forms of nuclear energy. Historical and other traveling exhibits are also displayed in this area. Located outside of the museum are a number of large exhibits. These include the Boeing B52B jet bomber and a Navy TA-7C Corsair II fighter-bomber as well as many other nuclear weapons systems, rockets, and missiles.

Robotic Manufacturing Science & Engineering Laboratory Tour

Intelligent systems bring diverse technologies together: computers, software, sensors, vision systems, and hardware such as robots. At Sandia National Laboratories, combinations of these technologies are merged to create robotic and intelligent systems that range from micro to mega.

To advance the evolution of robotic and intelligent system technologies, Sandia National Laboratories and the DOE created the Robotic Manufacturing Science and Engineering Laboratory (RMSEL). It is the first centralized facility designed specifically for bringing intelligent machine technologies and technologists together.

The RMSEL facility was designed as a special environment to accommodate the unique needs of robotics and intelligent systems research. A second-floor viewing gallery concourse overlooks ground-floor laboratories used for the development of large-scale robotics systems. The State-of-the-art physical resources coupled with outstanding intellectual resources make RMSEL unique in robotic and intelligent systems research and development.

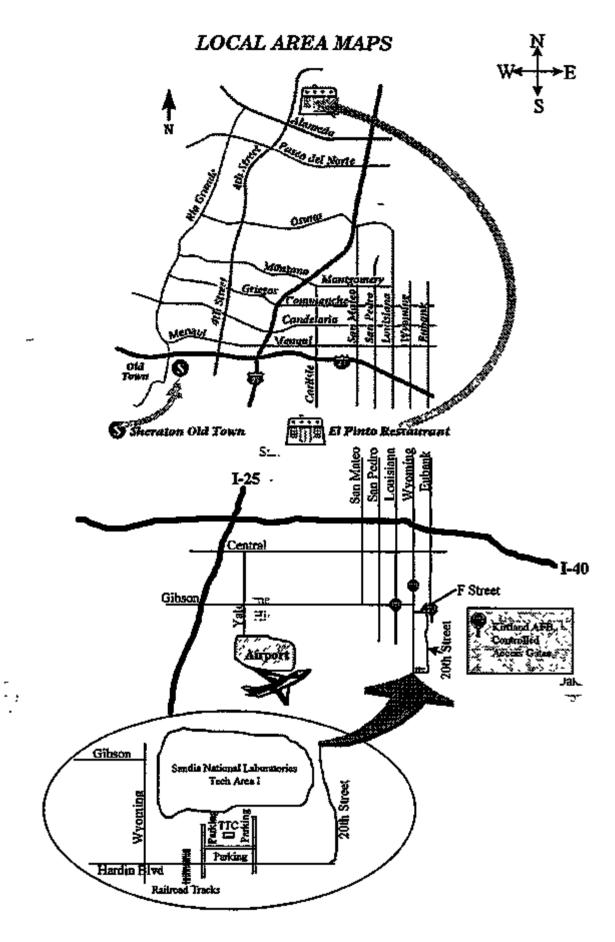
One of the main purposes of RMSEL is encouraging collaborative development with industry and academic partners.

No Host Dinner - El Pinto Restaurant

A No-Host dinner has been planned for Wednesday Evening at the El Pinto Authentic New Mexican Restaurant located at 10500 4th NW. There will be a variety of dinner selections offered that should accommodate all tastes. The cost of the dinner is \$15. Check at the Registration Desk in the TTC Lobby if you would like to attend or if you are planning to use the bus transportation provided from the Sheraton Hotel to the El Pinto Restaurant. El Pinto is located at 10500 4th NW; the phone number is 898-1771.

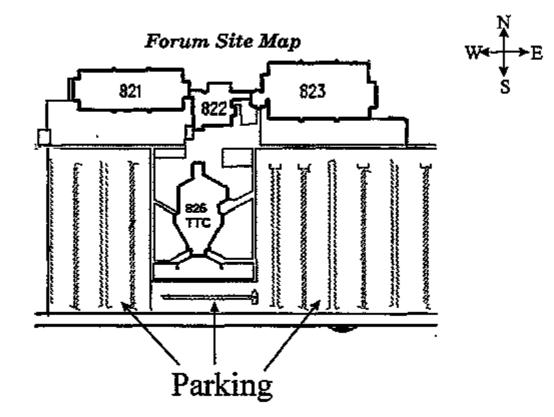
		lay, April 1, 1997			
Depart	Time	Destination	Return to Sandia National Labs, "Pick-Up" time		
Sandia National Labs, TTC	5:30 p.m.	National Atomic Museum (Social)	6:30 p.m.		
Veinesday, April 2, 1997					
Sandia National Labs, TTC	1:30 p.m.	Robotics Lab	2:45 p.m.		
Sandia National Labs, TTC	1:30 p.m.	National Atomic Museum	2:45 p.m.		
Sandia National Labs, TTC	3:15 p.m.	Robotics Lab	4:45 p.m.		
Sandia National Labs, TTC	3:15 p.m.	National Atomic Museum	4:45 p.m.		
Sheraton Old Town	5:45 p.m.	El Pinto Restaurant	Return to Sheraton Old Town, "Pick-Up" time ~8:30 p.m.		

Bus Schedule for Social, Tours, No-Host dinner



Ĵ

ſ



Location of Conference Rooms

TTC Auditorium, TTC Lobby, TTC Conference Room

Located in Building 825. Enter through the doors on the north side of the building.

Bldg. 822 Rooms A&B

Located immediately to the right when entering Bldg. 822 from the doors on the south side of the building.

Bidg. 823 Breezeway

Located immediately to the left after the Reception Desk when entering Bldg. 823 from the doors on the south side of the building.

NOTE: To get into the 823 Breezeway, individuals without a valid DOE must be escorted by an individual with a valid DOE badge. They must show a picture ID and sign in at the reception desk. The Breezeway will only be used for the afternoon Keynote tutorials and a Forum committee member will be available to assist you with the entrance details.

()



Biographies & Abstracts

:.

BIOGRAPHIES

۰.

2

r



Keynote Biographies

Capers Jones, Chair SPR

Capers Jones is an international consultant on software management topics and Chairman of Software Productivity Research, Inc. (SPR) in Burlington, MA. Following graduation from the University of Florida, Mr. Jones began his software career as a programmer in the office of the Surgeon General, Washington, D.C.. Prior to becoming Chairman at SPR, Mr. Jones also worked at the Crane Company, IBM, and was Assistant Director of Programming Technology at ITT in Stratford CT. Mr. Jones has published nine books dealing with software areas including; programming productivity, software measurement, and software quality. His tenth book, <u>Software Cost Estimating</u> is scheduled for publication in early 1997. Mr. Jones will share his experience and insights in his keynote address "Software Ouality for 1997 - What Works and What Doesn't".

Presentation: April 2, (09:00-10:00 am), TTC Auditorium



Dr. David Lorge Parnas, McMaster University

Professor David Lorge Parnas, Ph.D. holds the NSERC/Bell Industrial Research Chair in the Communications Research Laboratory, Department of Electrical and Computer Engineering at McMaster University in Hamilton, Ontario, Canada. His primary area of interest is to promote to Software Engineers the discipline and body of knowledge as practiced by engineers in other fields.

By studying the problems of software engineering since 1965, Dr. Parnas has developed principles and methods that have value to real world problems. In reservation of his accomplishments, he has received second honors, including election as a Fellow of



the Royal Society of Canada and a Fellow of the Association for Computing Machinery. • Dr. Parnas will share his experience and knowledge by leading three workshop/tutorials. Tutorials: April 1, ZI (09:00-11:00 am), TTC Auditorium Z2 (01:00-03:00 pm), Z3 (03:15-05:15 pm), Bldg 823 Breezeway

Tutorial Leader Biographies

(Alphabetical Order)

Nancy L. Crowley, Phillips Laboratory

Lt Col Nancy Crowley is the Acting Chief of the Space System Technologies Division (PL/VTS), Kirtland AFB, New Mexico. The focus of Space System Technologies Division is on the innovative application of software technologies to improve performance and reduce operations and maintenance costs for satellite control systems, including telemetry, tracking and commanding (TT&C), mission data dissemination, data processing, and satellite autonomy. Lt Col Crowley is also the program manager for the Multimission Advanced Ground Intelligent Control (MAGIC) program. MAGIC is developing the architecture for the next generation satellite control system that provides a low cost, flexible software architecture that allows plug and play of COTS products in a vendor independent manner. Lt Col Crowley was born May 13, 1955 in the Bronx, New York. She graduated from Theills High School in Theills NY, in 1973. She received a Bachelor of Science in Electrical Engineering from the University of New Hampshire in 1977 where she was a ROTC distinguished graduate. She later received the Master of Science in Digital Engineering and the Doctor of Philosophy (major of software engineering, minor of artificial intelligence) from the Air Force Institute of Technology in 1982 and 1994 respectively. Her research was in object-oriented methods for software requirements analysis. Lt Col Crowley entered the Air Force in 1972 and was a flight test engineer for Tactical Air Command. There she conducted operational test and evaluation and flew in fighter aircraft in support of projects. After her masters degree, she was assigned to the Flight Dynamics Laboratory, where she was the software engineer for the digital flight control system of the X-29 Advanced Technology Demonstrator and the Ada focal point for the laboratory. There and in subsequent assignments she was a technical consultant to the Swedish government on the development of the digital flight control system for the JAS-39. Her next assignment was at the Systems Acquisition School, Brooks AFB Texas where she was a course developer and instructor of software acquisition courses. There she was also a system administrator for a UNIX and PC-based networked system that serviced the students and staff at the school. After completing her Ph.D., she came to her current assignment in Oct 94. Outside her Air Force duties, Lt Col Crowley teaches software engineering, software management, and computer science courses at local Universities. Her and her busband own a computer consulting business. Both her and her husband enjoy riding horses. Tutorial X2: April 1, (03:15 - 05:15 pm), SNL Bldg 822, Room B

Randy Dabbs, Sandia National Laboratories

Randy Dabbs is a Senior Member of Technical Staff at Sandia National Laboratories. He has earned a Master of Science in Electrical Engineering from the University of New Mexico. He has held positions at the Sandia Particle Beam Fusion Accelerator in the areas of data acquisition and signal processing; the Kwajalein Missile Range in the areas of range computer systems engineering, range operations, tracking software modeling and development, reentry mission project engineering, digital radar signal processing, radar controller real time software, and software configuration management; and the Sandia Kanai Test Facility in the areas of range computer support and operations, range safety software development, countdown software development, CASE tool selection and modeling of range operational software. In his current position with the Sandia Quality Engineering Department, he has participated in instructing the Software Quality Engineering course and the Software Inspections course. In his role as software quality assurance engineer, he has participated in numerous software inspections for both internal and external customers. In addition, he has helped develop and teach a customized version of the software inspection course to meet the specific needs of Sandia organizations.

Tutorial Y2: April 1, (03:15 - 05:15 pm), TTC Conference Room C

Dwayne L. Knirk, Ph.D., Sandia National Laboratories

Dr. Knirk is a member of the software quality engineering department at Sandia National Laboratories. He provides in-house consulting to line organization projects for software engineering processes, methods, standards, tools, and training. He participates in process assessments and improvement programs, and provides support for configuration management, software inspections, and process automation. Dr. Knirk's primary focus is on the two complementary areas of software specification and testing, in which he works to bring more formal methods into more practical applications. He works actively on IEEE software engineering standards groups. He is a member of the ASQC Software Division Methods Committee. Dr. Knirk previously worked for Programming Environments, Inc., where he was the architect and principal developer of the automated software test design tool, T. That commercial product analyzed a formal software behavior description for testability, designed test cases for demonstrating that behavior, and generated actual test case data.

Tutorial W2: April 1, (03:15 - 05:15 pm), SNL Bldg 822, Room A

Tutorial Leader Biographies

(Alphabetical Order)

G. Lawrence Lane, Sandia National Laboratories

Larry Lane is a Senior Member of the Technical Staff at Sandia National Laboratories. He earned a Master of Arts Degree in mathematics from the University of Kansas. Larry joined Sandia Corporation in 1959 as an assembly language programmer in the field data reduction department. He has also worked as a operating systems programmer and was responsible for the selection and installation of Sandia's first general purpose time sharing computer. Larry also worked as a computer consultant for large scientific computers, as the second computer onbudsman, and was responsible for the development of an electronic tracking system for electrical testing of radiation-hardened microclicuits. Larry moved to his corrent position in the Quality Engineering Department in 1991, where he is an instructor for the Software Quality Engineering course and the Software Inspection Class. As a software quality engineer, Larry has led numerous qualification efforts for new and upgraded software projects, particularly in the areas of use control and weapon security. He has helped develop and teach a customized version of the software inspection course to meet specific Sandia organizational needs.

Tutorial Y2: April 1, (03:15 - 05:15 pm), TTC Conference Room C

Gerald W. McDonald, Ph.D.

Dr. McDonald has a Bachelor of Science in Engineering Science and a Master of Science in Computer Systems Management from the Naval Postgraduate School. Following his retirement the Navy he received a Master of Engineering in Industrial and Systems Engineering and a Ph.D. in Quantitative Management Science (Operations Research) from the University of Florida. Following receipt of his Ph.D. he worked for BDM International as an executive-level Program and/or Project Manager and technical leader. During his thirteen years with that firm he led both software and non-software projects. During the three years since his retirement from EDM he has acted as consultant to Sandia, SEMATECH, and a number of other organizations. As a consultant he has worked primarily in the field of Software Process Improvement. Besides direct technical assistance he has presented training and workshops in software areas such as: quality engineering, software inspections, process definition and documentation, and metrics.

Tutorial XI: April 1, (01:00 - 03:00 pm), SNL Bldg 822, Room B

John K. Sharp, Ph.D., Sandia National Laboratories

John has performed information analysis in various positions at Sandia for fifteen years. He has worked closely with Prof. Shir Nijssen of the Netherlands for several years to establish the procedure to develop and analyze information problems using structured natural language. They are currently finishing a text on this topic. This procedure was originally based on the NIAM (Natural language Information Analysis Methodology) modeling technique. John and Prof. Nijssen have co-chaired two international conferences on natural language modeling. John is also the editor of the international standard on conceptual schemas. *Tutorial W1: April 1, (01:00 - 03:00 pm), SNL Bldg 822, Room A*

Software Quality Assurance Subcommittee, Work Item #16, Nuclear Weapons Complex Sites The Software Quality Assurance Subcommittee (SQAS) operates under the DOE Nuclear Weapons Complex (NWC) Quality Managers to identify and resolve Software Quality issues and problems common to all DOE sites and facilities. This tutorial is the result of an NWC SQAS work item to define how to manage and control software as product. The work item was established to satisfy a need to define a consistent process for handling product software. The Nuclear Weapons Complex-wide participants and presenters of this tutorial include:

Chair David Vinson, Pantex Plant Phil Huffman, Pantex Plant Alvin Cowen, Pantex Plant Catherine Kuhn, AS/FM&T Donald Schilling, AS/FM&T Dave Peercy, SNL Mike Blackledge, SNL Orval Hart, LANL John Cerutti, LANL Bill Warren, LLNL Charles Chow, LLNL Ellis Sykes, DOE/Kansas City Area Office Gary Echert, DOE/Albuquerque Area Office Kathleen Canal, DOE/HQ Ray Collen, SRS Faye Brown, LMES, Oak Ridge, Y-12 Plant ð.

Tutorial Y1: April 1, (01:00 - 03:00 pm), TTC Conference Room C

(Alphabetical Order)

John Ambrosiano, Ph.D, Los Alamos National Laboratory

Dr. Ambrosiano received his Ph.D. in Plasma Physics from the College of William and Mary in 1980 and has since pursued a career in Computational Physics. He has written simulation codes for a variety of applications including plasmas and beams, acoustics, fluid dynamics, and electromagnetics. After a postdoctoral appointment at the University of Alaska's Geophysical Institute to study Space Physics, he moved to the Washington, DC area to work with a defense contractor. In 1987 he joined the Lawrence Livermore National Laboratory where he worked on nuclear weapons applications, and later joined the Earth System Modeling project there. The growing complexity of mimerical simulations led to a strong interest in Computer Science and in Software Engineering in order to find the leverage to manage the complexity of the new generation of simulation codes. In 1995 he joined the North Carolina Supercomputing Center to lead the effort to build a simulation framework for environmental modeling called the Environmental Decision Support System. This became the prototype for EPA's Models-3 framework. He recently joined Los Alamos National Laboratory to participate in DOE's Accelerated Strategic Computing Initiative. He is currently the leader of a twelve-person visualization and human-computer interaction team in X Division at LANL. He is also the Laboratory's principle investigator for Scientific Data Management within the ASCI program. His current interests are scientific data management, computational frameworks, and software engineering for scientific applications.

Presentation: Wednesday, April 2, Session C1: 10:15-11:45 am, TTC Conference Room C

Rodema Ashby, Sandia National Laboratories

Rodema Ashby has been programming or leading projects at Sandia for the last 13 years. Projects have included configurable software security systems such as the Site Independent Alarm and Display System, and a Logging and Accountability Subsystem. Interactive Collaborative Environments (ICE) which was licensed to SUN Microsystems as their "Show Me" product included a great deal of commercial customer testing and collaboration. A-PRIMED which was a 22 month, 2.5 million dollar cooperative effort involving 10 SNL NM Centers (and minimally KC and SNL CA), demonstrated a 24 day, new product to market cycle. New hardware from new customer requirements was created in a matter of days, after the project realization team had set up a communications network and created and integrated tools for product realization. Rodema is currently writing code to customize solid modeling tools for easier user model modifications.

Presentation: Wednesday, April 2, Session A1: 10:15-11:45 am, TTC Auditorium

Mikhail Auguston, New Mexico State University

Received a Ph.D. degree from the Institute of Cybernetics in Kiev (USSR) in 1983, Diploma of the Senior Research Fellow from the Highest Evaluation Commission of the Council of Ministers of USSR in 1990, and degree of Doctor in Computer Science from University of Latvia in 1992. Research interests are in programming language design and implementation, and program testing and debugging tool design.

Joined Computing Ceater of Latvia University as Research Scientist in 1971. Since 1983 worked as a Leading Researcher at the Institute of Mathematics and Computer Science of Latvia University. Took part in the design and implementation of the language for file processing, the interpreter for PL/1 program testing, the testbed environment for assembler level language for PDP-11 computers, the implementation of specification language SDL for communication system software rapid prototyping and testing, the tool system GRAPES/4GL for information system design. In the years 1987-88 has designed and implemented programming language RIGAL for compiler writing on PDP, VAX and IBM PC computers. This work was presented at a number of international conferences and is used at several sites for language processor design. In 1990 he has started to work on program formal annotation language FORMAN for sequential and parallel program dynamic analysis, testing and debugging. This work was presented at a various international conferences and in several universities in Europe and United States as an invited talk. He is the author of more than 30 scientific articles and co-author of the most popular textbook on PL/1 in Soviet Union (totally more than 100,000 copies printed). Currently he is an Associate Professor at the Computer Science Department of New Mexico State University. He teaches undergraduate and graduate classes on C++, Data Structures, Software Engineering, Compiler Construction, Ada programming language. Member of ACM and IEEE Computer Society.

Presentation: Thursday, April 3, Session B5: 10:15-11:45 am, Bldg 822, Rooms A&B

(Alphabetical Order)

Michael Bell, Lockheed Martin Energy Systems

Michael Bell is a software engineer with Lockheed Martin Energy Systems at the Y-12 Plant. He is the lead analyst on the Electronic Medical Records System project, as well as member of the software metrics team. He has worked in the Oak Ridge area for seventeen years, at both Y-12 and Oak Ridge National Laboratory. His experience includes research- and production-oriented software, in areas such as plasma physics, econometrics, access control, manufacturing, and inspection. In this capacity, he has performed user interface and database design, application migration (cross-platform and mainframe-to-workstation), real-time device control, modeling, statistical and graphical analysis, and all aspects of structured and object-oriented software development. Mike holds a bachelor's degree in mathematics and is currently working toward a master's degree in-software engineering.

i

Presentation: Wednesday, April 2, Session A2: 01:30-03:00 pm, TTC Auditorium

Gail M. Benefield, Lockbeed Martin Energy Systems

Ms. Benefield has worked for Lockheed Martin Energy Systems, Inc. (LMES) since 1987. Her assignments include working as an applications developer/analyst at the Y-12 site, an Applications Security Specialist for the Computing and Telecommunications Security Organization, and currently, as a Computing Specialist within the Information Technology Services division at the K-25 site in Oak Ridge. At Y-12, Ms. Benefield was on the team which revised the 80-Series, a document owned by the Y-12 Quality Division, which was the Y-12 implementation of the required software development methodology. She was also a member of the Y-12 Software Configuration Control Board, which reviews all software changes to applications which fall within a certain class of software. In her current assignment. Ms. Benefield is representing her department as an active participant on the team which authored and is supporting the Software WorkPackage Methods (SWM) methodology.

Presentation: Thursday, April 3, Session A4: 08:30-10:00 am, TTC Auditorium

Larry J. Dalton, Sandia National Laboratories

Larry J. Dalton holds a BS in Applied Mathematics and an MS in Electrical Engineering both from the University of New Mexico. Larry has spent the past 19 years at Sandia National Laboratories in Albumerove, New Mexico engaged in high consequence systems development. Much of that time was dedicated to various aspects of nuclear weapons and associated control systems. He is the manager of the Command and Control Software Department at Sandia National Laboratories which in addition to software engineering research, develops software and systems safety solutions for high consequence operations.

Presentation: Wednesday, April 2, Session B3: 03:15-04:45 pm, Bldg 822, Rooms A&B

Larry Desonier, Sandia National Laboratories

Education: In 1972, Larry graduated from Southwestern Louisiana with a Bachelors of Science in Electrical Engineering. In 1976 graduated from Oklahoma City University with a Masters in Business Administration. In 1979 completed Masters in Electrical Engineering and Computer Science from University of New Mexico. Complete a Masters of Science in Computer Information Systems from the University of Phoenix in 1996. Presently working on a Certificate in Computational Simulation Science from the University of New Mexico under a special Sandia National Laboratories retraining program with completion in May 1998. Work Experience: Officer in the U.S. Air Force from 1972 through 1975 and worked as a Communications-Electronics Engineer. Worked at the U.S. Air Force Weapons Laboratory from 1976 to 1984 as the Director of Communications. Came to Sandia National Laboratories in 1985 and has worked as a Systems Developer, Software Engineer, and Project Leader for over 12 years. -----

Presentation: Thursday, April 3, Session A5: 10:15-11:45 am, TTC Auditorium

John Hare, Ph.D., AWE, Ministry of Defence, United Kingdom

Dr John T Hare is the Software Quality Manager of AWE Aldermaston, an MOD (UK) facility managed by Hunting-BRAE Ltd. He is a Chartered Engineer and a Member of both the British Computer Society and the Institute of Quality Assurance. John graduated from the Universities of Nottingham (BSc) and York (DPhil). He started his career in 1973 as a scientist at what was then the Royal Aircraft Establishment (of International Airshow fame). He was responsible for analysis of sonobuoy trials data, using computers in the days when 16KByte was a generous amount of core memory! In 1980 John joined AEA Technology, which as UKAEA had been

(Alphabetical Order)

responsible for the UK Atomic Energy Programme. John was responsible for the design of a number of computerbased data acquisition systems. As the PC took the skill out of this activity, John's team specialised in Management Information Systems, and the provision of Software Engineering support to scientific projects. This was the start of a growing interest in Quality Assurance, as customers and regulatory authorities demanded accreditation to ISO9001. In 1993 John joined AWE, with a brief to improve software quality assurance and raise standards across the company. This is moving into a new phase, with emphasis on Software Engineering. John and his wife Heather have two daughters; Katherine (22) who is a biochemist doing research at Birmingham University, and Louisa (19) who is a student of Modern Languages at Nottingham University. Outside interests include local government and local history. Until recently John was Chairman of Governors at a school with 1000 students. *Presentation: Thursday, April 3, Session B5: 10:15-11:45 am, Bldg 822, Rooms A&B*

David L. Harris, Sandia National Laboratories

Dave has a M.S. in Computer Science and A.B in Mathematics from all from the University of Missouri. He was a graduate fellow at the Health Services Research Center in Columbia Missouri and his graduate education focused on multi-processor hardware architectures and multi-processing operating systems. Dave is currently a Senior Member of the Technical Staff at Sandia National Laboratories and is assigned to the Information Systems Engineering Center. Dave has been doing research in using World Wide Web technology in support of collaborative environments for distributed Decision Support Systems. Dave was the software process engineer for the ICADS (Integration Correlation and Display System) program. ICADS is a ground based satellite data analysis system and the project leader for TCAMS (Tech Control Automation, Maintenance, and Support), a five year, \$6 - 8M project consisting of over one million lines of software source code. TCAMS has been accepted by the Department of Defense customer and is in operation today. (A fielded and functional system). As the TCAMS Team Leader, Dave was responsible for the device control software subsystem of the TCAMS software project. Earlier in Dave's career he was a software engineer responsible for various systems analysis and design of a large command and control software system. Dave has software engineering experience in real-time, embedded, guidance and control computers for ballistic missiles and systems administration of large, multi-user, time-sharing systems.

Presentation: Wednesday, April 2, Session A1: 10:15-11:45 am, TTC Auditorium

Orval Hart, Los Alamos National Laboratory

Orval Hart has worked at the Los Alamos National Laboratory for 20 years, mainly involved in real-time control systems for nuclear facilities. He has a Bachelor's Degree in Mathematics from California State Polytechnic College (Cal Poly) at Pomona and a Master's Degree in Computer Engineering from the University of New Mexico. Prior to coming to Los Alamos, he worked in real-time data acquisition systems, later moving to the let Propulsion Laboratory in Pasadena where he worked on real-time telemetry and communication systems. In 1975, he moved to Los Alamos where he was responsible for the original building control system software for the Plutonium Research and Development facility (known as TA-55). Since then, he has worked on a control system for an unmanned nuclear power supply (later canceled), the original procurement of the Laboratory intrusion and detection system, an environmental monitoring computer network system for the Nevada Test Site and surrounding states, the facility control system for the Special Nuclear Materials Laboratory (a sister facility to TA-55 that was later canceled also), and for the last ten years has been responsible for the control software for the Weapons Engineering Tritium Facility. This system is not only a facility environment control system, but also assists in performing the everyday work in the Facility. Almost all work in the Facility is done from the control console as opposed to hands-on in glove boxes. As many of the procedural interlocks as could be foreseen were implemented in software to avoid human error, taking special care to test and prove them prior to going 'on-line'. Computer controlled automatic sub-systems are monitoring the Facility constantly to mitigate any operational abnormalities. This system was implemented during the early days of Admiral Watkin's tenure and as such, was a test case for increased compliance and formality-of-operations.

Presentation: Wednesday, April 2, Session C1: 10:15-11:45 am, TTC Auditorium

(Alphabetical Order)

Kevin Hill, Mason and Hanger Corporation, Pantex Plant

Kevin Hill is a tester design engineer at the Mason and Hanger Corporation. He holds a BS in electrical engineering from Kansas State University and is currently enrolled in the Interdisciplinary Master of Engineering curriculum at Texas Tech University. Co-author Dr. Mario G. Benvides is an assistant professor in Industrial Engineering at Texas Tech University. Dr. Beruvides has 10 years of industrial work experience in design, production, and manufacturing. His interests include white-collar/knowledge work performance improvement, productivity engineering, work measurement, technology management, and engineering education. Dr. Beruvides is a member of ASEM, a senior member of IIE, and a member of ASQC and the Academy of Management. He holds a BS in mechanical engineering and an MSIE degree from the University of Miami, and a Ph. D. from Virginia Polytechnic Institute and State University in industrial and systems engineering. *Presentation: Wednesday, April 2, Session B2: 01:30-03:00 pm, Bldg 822, Rooms A&B*

20

כווהנוס: יופפי

Curtis G. Holmes, Jr., Lockheed Martin Energy Systems

Curt came to Lockheed Martin Energy Systems (LMES) at Oak Ridge, Tennessee from Texas Instruments and is currently the Department Manager of the Environmental, Waste Management, and Analytical Laboratories Systems in the Data Research and Development Organization. The purpose of the department is to be a focal point for providing computing support for the Environmental, Waste, and Analytical Laboratory business areas at LMES. Prior to his current assignment, Curt was the Department Manager for the Computer Application's Department in the Engineering Division. The main focus of this department is the design, development, implementation, and deployment of digital systems to support real time process control and data acquisition systems. Curt Holmes holds a B.S. and M.S. Degree in Electrical Engineering from the University of Tennessee with a Minor in Computer Science. He is a licensed Professional Engineer in the State of Tennessee. *Presentation: Thursday, April 3, Session A5; 10:15-11:45 an, TTC Auditorium*

Karen Jefferson, Sandia National Laboratories in California

Karen L. Jefferson has worked at the Sandia National Laboratories for 12 years and is currently in the Systems Research Department at Sandia California. Her work experience at Sandia has included high performance computing, realtime control, software engineering, and systems analysis. She is currently the software project lead on the Advanced Atmospheric Research Equipment project. She has a Masters degree in Computer Science from the University of Arizona.

Presentation: Wednesday, April 2, Session A2: 01:30-03:00 pm, TTC Auditorium

Bruce L. Johnston, Mason & Hanger Corporation, Pantex Plant

Bruce L. Johnston is a Project Programmer/Analyst for Mason & Hanger Corporation at the DOE Pantex Plant. In April 1996, he accepted the challenge to be the Project Manager for the year 2000 Project. Before accepting this new assignment he was the Computer Security Site Manager for the Pantex Plant and has worked in a computer security capacity for the last ten years. Prior to joining Mason & Hanger, he worked for Battelle Memorial Institute in Richland, Washington, and with EG&G in Idaho Falls, Idaho. In his personal life he has served as a Scontmaster for his community and is currently serving as a Bishop for the Church of Jesus Christ of Latter-Day Saints. He keeps a healthy perspective and stays in balance by being a father of four children. *Presentation: Thursday, April 3, Session AS: 10:15-11:45 am, TTC Auditorium*

Marie-Elena C. Kidd, Sandia National Laboratories

Marie-Elena C. Kidd is a computer scientist and Senior Member of the Technical Staff at Sandia-National Laboratories. During her ten years at Sandia, she has worked as a software engineer on embedded, real-time software systems for such applications as robotics, nuclear weapon components, and control systems. She has also worked on lab-wide information sharing software systems and software engineering initiatives. She has a B.S. in Computing and Information Sciences, Trinity University, San Antonio, TX and an M.S. in Computer Science, Purdue University, West Lafayette, IN.

Presentation: Thursday, April 3, Session B4: 08:30-10:00 am, Bldg 822, Rooms A&B

(Alphabetical Order)

Dr. Dwayne L. Knirk, Ph.D., Sandia National Laboratories

Dr. Knirk is a member of the software quality engineering department at Sandia National Laboratories. He provides in-house consulting to line organization projects for software engineering processes, methods, standards, tools, and training. He participates in process assessments and improvement programs, and provides support for configuration management, software inspections, and process automation. Dr. Knirk's primary focus is on the two complementary areas of software specification and testing, in which he works to bring more formal methods into more practical applications. He works actively on IEEE software engineering standards groups. He is a member of the ASQC Software Division Methods Committee. Dr. Knirk previously worked for Programming Environments, Inc., where he was the architect and principal developer of the automated software test design tool, T. That commercial product analyzed a formal software behavior description for testability, designed test cases for demonstrating that behavior, and generated actual test case data.

Presentation: Wednesday, April 2, Session B1: 10:15-11:45 am, Bldg 822, Rooms A&B

Catherine M. Kubn, AS/FM&T Kansas City Site

Cathy Kuhn is a Staff Technical Programmer/Analyst from AlliedSignal Federal Manufacturing and Technologies / Kansas City Site. For the past eight years she has been a member of the Kansas City's Software Quality Assurance Group. During that time she has been involved in many Kansas City site and corporate software development and software quality improvement efforts. Currently, she is an active member of the Information Systems' Software Process Group and the Information Systems Software Quality Assurance Group. This presentation is based upon her work with the Information Systems' organization.

Presentation: Thursday, April 3, Session A4: 08:30-10:00 am, TTC Auditorium

Michael F. Lackner, AS/FM&T Kansas City Site

Michael holds a Masters of Science degree in Mechanical Engineering from the University of Missouri-Rolla, and a Bachelor of Science degree in Aerospace Engineering from the same institution. Michael is a Registered Professional Engineer in the State of Missouri. He is currently enrolled in the Doctor of Engineering program at the University of Kansas, specializing in the area of computer-aided and computer-integrated manufacturing. Prior to the SQA assignment eight years ago, he spent 4 years in process and product engineering in plastics products at AlliedSignal. He most recently completed the Blackbelt training in Six Sigma. Presentation: Thursday, April 3, Session B5: 10:15-11:45 am, Bidg 822, Rooms A&B

David J. Leong, Sandia National Laboratories

David has been a Senior Member of Technical Staff at Sandia National Laboratories for seven years. He is currently the project leader of Sandia's Internal Web Technology Team, the EVE (Enterprise-information Viewing Environment) Team. He has been involved with Sandia's Intranet from its inception in the summer of 1994. David has performed many related activities along the way, including; HTML authoring, browser training, systems integration, application development, browser/server installations, etc.. Sandia's Intranet, which has been featured in WebMaster Magazine and Netscape's Customer Profiles, currently honses approximately 40,000 administrative and technical documents and is accessed on the order of 250,000 times per day.

Presentation: Wednesday, April 2, Session B2: 01:30-03:00 pm, Bldg 822, Rooms A&B

Stewart Meyer, Savannah River Site

Stewart Meyer is currently the software Quality Assurance/Configuration Management Coordinator for the NWPS (Nuclear Waste Processing Support) section for all systems supporting the DWPF (Defense Waste Processing Facility) at SRS (Savannah River Site.) This position involves developing/updating QA/CM plans for process control, process support, and manufacturing support systems. He also performs a hands on role as the configuration manager for the SCMS (Software Configuration Management System) in developing the layered applications, reviewing and approving the software changes, and performing library maintenance. He is the lead for all external (DOE/Site) audits regarding software at DWPF and also participates in committees and task teams at the division and Site level regarding software management procedures. A graduate of McMurry College (Abilene, Texas), with a Bachelor of Science in Computer Science and a background in management, his software engineering career includes; OS/Application development for the DOD MLRS (Multiple Launch Rocket System)

(Alphabetical Order)

project, process automation design/development for DWPF, group supervisor for the process automation group at DWPF, and his current position (since 1993.)

Presentation: Wednesday, April 2, Session A2: 01:30-03:00 pm, TTC Auditorium

Jennie L. Negin, Sandia National Laboratories

Jennie Negin is manager of Web Services and IS Training at Sandia National Laboratories in Albuquerque, New Mexico. Sandia is a Department of Energy multiprogram national laboratory managed by Sandia Corporation, a Lockheed Martin company. Ms. Negin has been involved in development of many Information Systems at Sandia - travel, library, procurement, property, security, personnel, nuclear materials management and radiation exposure. Ms. Negin was a consultant to the University of New Mexico (UNM) Law School and the UNM Maxwell Museum of Anthropology before coming to Sandia. Prior to that she was an internal consultant and systems developer at Los Alamos National Laboratories and the University of Florida Computing Center. Ms. Negin is a long time member of the Association of Computing Machinery and the New Mexico Network for Women in Science and Engineering. Jennie is a graduate of the University of Florida with a BSE and MA in Mathematics. *Presentation: Wednesday, April 2, Session B2: 01:30-03:00 pm, Bidg 822, Rooms A&B*

Don Rathbun, AS/FM&T Kansas City Site

Don Rathbun holds a BSEE from Kansas State University, Manhattan, Kansas, and a MSEE from the University of Missouri, Columbia, Missouri. Business Systems Reengineering has been the focus of Don's recent assignments including project responsibilities on the Focused Factory initiative and the ISO9001 certification process from its outset. Current assignments include involvement with the NWIG (Nuclear Weapons Information Group), IMOG (Interagency Manufacturing Operations Group), and CAM-I (Consortium for Advanced Manufacturing International) Organizations. Don has made presentations at the last two IMOG meetings and at the September 1995 LLNL Software Engineering Seminar. Prior assignments included project responsibilities on major radar fuzing systems.

Presentation: Wednesday, April 2, Session A3: 03:15-04:45 pm, TTC Auditorium

Larry Rodin, Mason & Hanger Corporation, Pantex Plant

Lasty has been 30 Years with Mason & Hanger Corporation working in Quality. He is a Project Manager at the Pantex Plant, Amarillo, Texas, Senior Member of the American Society for Quality Control, Member Software Quality Division. Larry has been an ASQC Certified Quality Engineer since 1970. In deference to the Year 2000 phenomena, his recertification date is December 31, 1999. Larry became Mason & Hanger's SQAS Primary Representative in the fall of 1990. He is currently serving as SQAS Vice-Chair, and previously has served as Secretary. Larry has also worked on many Work Item Groups and developed this presentation as research for one of these groups.

ł

.

1

ł

ţ.

Presentation: Thursday, April 3, Session B5: 10:15-11:45 am, Bldg 822, Rooms A&B

Edward W. Russell, Lawrence Livermore National Laboratory

For the last 15 years Ed Russell has been involved in formal QA implementation on several projects at LLNL. He is currently working toward the ASME NQA-1 lead auditor qualification. Ed has also worked as an FEM code analyst at LLNL in the early 1980's. Ed's academic achievements include an M.S. degree from the University of California Davis in Mechanical Engineering and Materials Science.

Presentation: Wednesday, April 2, Session CI: 10:15-11:45 am, TTC Auditorium

Don Schilling, AS/FM&T Kansas City Site

Don Schilling is a Manager, Engineering Projects, for AlliedSignal Federal Manufacturing and Technologies at Kansas City. He has over 30 years of manufacturing experience in various assignments and responsibilities. He was responsible for the formation of the Kansas City Plant's Software Quality Assurance Group, which has reported to him since 1988. Don has championed numerous Software Engineering and SQA initiatives within AlliedSignal, the DOE Nuclear Weapons Complex, and in national and international forums. *Presentation: Wednesday, April 2, Session A3: 03:15-04:45 pm, TTC Auditorium*

(Alphabetical Order)

Joseph R. Schofield Jr., CQA, Sandia National Laboratories

Joe has been applying emerging technology for business and engineering solutions for the past 17 years. Joe guided the evaluation and implementation of Sandia's first large-scale CASE project using Texas Instrument's IEF. Current efforts include a client-served based object-oriented project with tens of millions of object instances. Joe has been a keynote speaker at the Structured Development Forum in San Francisco in 1988 and spoke on CASE at the National Conference on Information Systems Quality Assurance in Orlando, CASEWorld in LA, and the Pledmont CASE User's Group in Charlotte. Several articles on CASE were published by the Journal of Quality Data Processing, System Builder, and Managing System Development. A four-page interview was printed in the CASE Strategies Newsletter and another in Government Computer News. Joe has presented at USE, SHARE, GUIDE, and DOE-sponsored conferences. The Next Silver Buillet was published in 1995. His most recent article The Year 2000 - Finally a Reality Check is under publication review.

Presentation: Wednesday, April 2, Session A1: 10:15-11:45 am, TTC Auditorium

John K. Sharp, Ph.D., Sandia National Laboratories

John has been working in information systems during a 16 year career at Sandia National Laboratories. He has held technical and management positions covering information system design, application development and data administration functions. In has been working closely with Professor Shir Nijssen in the Netherlands who is creator of the NIAM (Nijssen's Information Analysis Methodology), which is the basis for our approach to Natural Language Modeling. Shir and John have co-chaired two international conferences on Natural Language Modeling and are writing a book on Natural Language Modeling that will be published this winter.

Presentation: Thursday, April 3, Session B5: 10:15-11:45 am, Bldg 822, Rooms A&B

Debra Sparkman, Los Alamos National Laboratory

Debra Sparkman is the Software Quality Assurance Manager for LLNL Safeguards and Security Engineering and Computations Division. She has been the SSEC quality assurance manager since January 1993 and test coordinator for the Argus Security System since October 1994. Prior positions at LLNL have included Quality Assurance/Test Coordinator for the Controlled Material Tracking System and staff member for the Fission Energy and Systems Safety Computer Safety and Reliability group. Other publications include: SSEC SEI Experiences, 1994 DOE NWC Software Quality Forum and Standards and Practices for Reliable Safety-Related Software Systems, 3rd International Symposium on Software Reliability Engineering. Ms. Sparkman received a Bachelor of Science, Computer Science in 1984 from the University of the Pacific. She is a member of the American Society for Quality Control, IEEE, and IEEE Computer Society.

Presentation: Wednesday, April 2, Session B1: 10:15-11:45 am, Bldg 822 Rooms A&B

Ann Stewart, Lockheed Martin Energy Systems

Ms. Stewart is the Quality Manager of the Data Systems Research and Development Program (DSRD) a division of Lockheed Martin Energy Systems (LMES) in Oak Ridge, Tennessee. She has more than 20 years experience as a software engineer and project manager with extensive experience in areas of quality assurance, performance measurements, and process improvement. She established and managed the Software Quality Assurance Program for the Oak Ridge National Laboratory (ORNL) in compliance with the Department of Energy (DOE) requirements and was responsible for their Performance Indicator and Metrics Program. Ann is a graduate of the University of Teanessee with a B.S. in Computer Science. She currently leads and manages DSRD's Process Improvement Initiative using the Software Engineering Institute's Capability Maturity Model (SEI/CMM). *Presentation: Thursday, April 3, Session A4: 08:30-10:00 am, TTC Auditorium*

Nancy A. Storch, Lawrence Livermore National Laboratory

Nancy has over 30 years experience in design and development of scientific software, with emphasis in user interface design, computer graphics and software engineering. Her special interest is usability engineering. Recently Nancy has also become involved in software quality assurance and serves as SQA Engineer to two projects. Nancy is the LLNL SE/SQA Group Leader. Prior to coming to LLNL, Nancy developed software for submarine fire control systems. Throughout her career, Nancy has striven to be at the forefront of the application of computer science and software engineering. She has done graduate work in human factors, user interface design, computer science and physics. Her degree is in mathematics.

Presentation: Wednesday, April 2, Session B1: 10:15-11:45 am, Bldg 822 Rooms A&B

(Alphabetical Order)

Michael Tiemann, Headquarters Department of Energy

Mike Tiemann has served in government service for 25 years. His career started in 1972 at Army Material Command Headquarters, as an Army Lieutenant working in Environmental Program Management. After this he spent almost 13 years at the Federal Energy Regulatory Commission as an Environmental Protection Specialist and a Computer Systems Analyst. In 1987 he joined Headquarters DOE as the Project Management Officer coordinating all information technology services and support for the Offices of the General Council, Inspector General, Hearings and Appeals and the Economic Regulatory Administration and the Board of Contract Appeals. Two years later, he was assigned the primary responsibilities for Information Management Planning at Headquartera. He is currently the Action Officer in the CIO's Information Architecture Team responsible for development of the Departmental Information Architecture. He is also the leader of the Information Management ' Planning and Architecture Coordinating Team or IMPACT, a diverse and professionally robust group of technology professionals from across the Department which supports the Architecture efforts. In addition to IMPACT, Mike has been a member of several Department-wide teams, and recently sat on an interagency panel on business modernization. Mike holds degrees in Architecture (BED, Texas A&M, 1972) and Systems Management (MSSM, U.S.C., 1977). He is a current member of the Energy Federal Credit Union's Information Technology Advisory Committee. He is married and has two children.

ļ

t

ì

ţ

5

Т

L

i

÷

i

Presentation: Wednesday, April 2, Session A3: 03:15-04:45 pm, TTC Auditorium

Victor L. Winter, Ph.D., Sandia National Laboratories

Victor L. Winter received his Ph.D. from the University of New Mexico in 1994. His dissertation research focused on proving the correctness of program transformations. Currently, Dr. Winter is a member of the High Integrity Software (HIS) Project at Sandia National Laboratories. His research interests include trusted software, formal semantic models (graphical-based and symbol-based), theory of computation, automated reasoning and robotics. Dr. Winter can be reached by phone in the United States at (505) 284-2696, by fax at (505) 844 - 9478, or by email at *viwinte@sandia.gov*.

Presentation: Wednesday, April 2, Session B3: 03:15-04:45 pm, Bldg 822, Rooms A&B

Alexander R. Yakhnis, Ph.D., Pioneer Technologies

Dr. Alexander R. Yakhnis is a consultant in design of dependable software/hardware systems. He received a Diploma in Mathematics from Moscow State University, Moscow, Russia. He worked as a computer programmer in Moscow, Russia and Houston, Texas. Alexander received an M.S. in Computer Science and a Ph.D. in Mathematics/Computer Science from Cornell University, Ithaca, New York. He then worked as a Research Scientist at Maihematical Sciences Institute, Cornell University. He worked at Command and Control Software Department at Sandia National Laboratories on High Integrity Software project from July 1995 to August 1996. His interests include correctness proofs for concurrent and sequential programs, theory of computations, winning strategies for two person games, control theory, hybrid systems, object-oriented methods, design of hardware/software systems. He can be reached by phone at (505) 298-5854 or by e-mail at AYakhnis@aol.com. Co-suthor Dr. Vladimir R. Yakhnis is a research scientist at Rockwell Science Center, One Thousand Oaks, CA. He received a Diploma in Mathematics from Moscow State University, Moscow, Russia. He worked as a computer programmer in Moscow, Russia and Houston, Texas. Dr. Yakhnis received an M.S. in Computer Science and a Ph.D. in Mathematics/Computer Science from Cornell University, Ithaca, New York. His research was in program correctness for concurrent and sequential programs, winning strategies for two person games, state transition systems and object-oriented methods. Dr. Yakhnis worked at the IBM Endicott Programming Laboratory as an Advisory Programmer until 1994. There he developed "Generic Algorithms" methodology that allowed the construction of mathematically proved software while "hiding" the actual proofs from the developers. The methodology was designed to take advantage of object class templates in C++ or Eiffel. He worked as a Visiting Scientist at Mathematical Sciences Institute, Cornell University until June 1995. There he developed the groundwork for the semantics of object-oriented stepwise refinements. He worked at Sandia National Laboratories at Albuquerque during 1995-1996. He can be reached by phone at (805) 373-4856 or by e-mail at vryakhni@scimail.risc.rockwell.com.

Presentation: Wednesday, April 2, Session B3: 03:15-04:45 pm, Bldg 822, Rooms A&B

ABSTRACTS



.

Tutorial Abstracts: Tuesday, April 1 1997

Keynote Tutorial 09:00 - 11:00 am

Dr. David Lorge Parnas, MU/CRL

20: Design Through Documentation: The Path to Software Quality

TTC Auditorium

Although it is appealing, practitioners are not able or willing to write precise documents. Instead, they write vague blarbs that are useless to those charged with the next steps and cannot be subject to rigorous analysis. This tutorial describes how precise, complete, and testable documents can be produced for software and the ways that these documents can contribute to an improved software process.

Tutorials 01:00 - 03:00 pm

Dr. David Lorge Parnas, MU/CRL Z1: Inspection of Critical Software

Plds 022 Deserves

Bldg.823 Breezeway

This tutorial describes a procedure for inspecting software that consistently finds subtle errors in "mature" software, software that is believed to be correct. The procedure is based on three key ideas: the software reviewers are active not passive; reviewers focus on small sections of code; reviewers proceed systematically so that no case and no section of the program gets overlooked. During the procedure, the inspectors produce and review mathematical documentation. The mathematics and its notation allows them to check for complete coverage and to proceed systematically and in small steps.

Dr. John Sharp, Sandia National Laboratories

WI: Natural Language Modeling

Bidg 822 Room A

This tratorial describes a process and methodology that uses structured natural language to enable the construction of precise information requirements directly from users, experts, and managers. The main focus of this natural language approach is to create the precise information requirements and to do it in such a way that the business and technical exerts are fully accountable for the results.

Dr. Gerald McDouald, Sandia National Laboratories Consultant

X1: Definition and Documentation of Engineering Processes Bldg \$22 Room B

This tutorial is an extract of a two-day workshop developed under the auspices of the Quality Engineering Department at Sandia National Laboratories. The presentation starts with basic definitions and addresses why processes should be defined and documented. It covers three primary topics: (1) process considerations and rationale, (2) approach to defining and documenting engineering processes, and (3) an IDEFO model of the process for defining engineering processes; describe the general architecture for product development; and define implications of immature processes versus those for mature processes. The approach describes the top-level subprocesses that make up the methodology for definition and documentation of engineering processes; namely: planning, gaining management approval for a process definition project, collecting data on the as-is process to capture current best practices within the organization, constructing a model of the as-is process, and verifying and validating that model. The final portion presents a four-level, hierarchical model that describes HOW to define and document an engineering processe.

I.

ļ

Fay: Brown, Oak Ridge; Ray Collen, Savannah River; Gary Echert, DOE/AL; Phil Huffman, Panter. Cathy Kuisa, AS/FM&T; Dave Peercy, SNL; Ellis Sykes, DOE/KCP; David Vinson, Panter VI: Now the NWC Handles Software as a Product

TTC Conference Room C

The statistical provides a hands-on view of how the Nuclear Weapons Complex projects should be handling software as a product in response to Engineering Procedure 401099. The primary scope of the interial is on software products that result from weapons and weapons-related projects, although the information presented is applicable

Tutorial Abstracts: Tuesday, April 1 1997

to other software projects. Processes for Identification, Qualification, Acceptance, and Delivery are described in terms of an extended case study.

Participant Restrictions: Must be a NWC or government employee; identification will be required. If you have questions, contact Dave Peercy, 505-844-7965, depeerc@sandia.gov.

Tutorials 03:15 - 05:15 pm

Dr. David Lorge Parnes, MU/CRL Z2: Exercise and Discussion Bide 273 Broomerson

Bidg 823 Breezeway

In this workshop, participants will be given a small program and will apply the documentation and inspection's methods from the previous Design Through Documentation and Inspection of Critical Software tutorials. This will be followed by a discussion of previous experiences in a question and answer format.

Participant Restrictions: Must have attended both the Design Through Documentation and Inspection of Critical Software tutorials.

Dr. Dwayne Knirk, SNL

W2: Writing Testable Software Requirements

Bidg 822 Room A

This tutorial identifies common problems in analyzing requirements in the problem and constructing a written specification of what the software is to do. It deals with two main problem areas: separating the documentation of what is given from the documentation of what is to be created; and determining what facts about the subject software are to be documented, how they should be expressed, and how they are related.

Lt. Col. Nancy Crowley, USAF Phillips Laboratory

X2: Using COTS Software in Development Projects

Bidg 822 Room B

Commercial software and standards must be carefully evaluated prior to selection, carefully integrated, and used where appropriate to reap their benefits. This tutorial will discuss the experiences of the Space System Technologies Division of the USAF Phillips Laboratory in developing a COTS-based satellite control system.

Larry lane and Randy Dabbs, Sandia National Laboratories

Y2: Software Inspection Process Overview

TTC Conference Room C

This tutorial provides an overview of the Software Inspection (In-Process Formal Review) Process and a miniinspection workshop. The inspection roles and process steps are introduced. Participants are then divided into inspection groups for conduct of a mini-inspection to gain some practical experience with the inspection process. Discussion of the mini-inspection results concludes the workshop.

Keynote Address, 09:00 - 10:00 am, TTC Auditorium

Capses Jones, McMaster University

Softx ere Quality for 1997 - What Works and What Doesn'1?

This presentation provides a view of software quality for 1997 – what works and what doesn't. For many years, software quality assurance lagged behind hardware quality assurance in terms of methods, metrics, and successful results. New approaches such as Quality Function Deployment (QFD) the ISO 9000-9004 standards, the SEI maturity levels, and Total Quality Management (TQM) are starting to attract wide attention, and in some cases to bring software quality levels up to a parity with manufacturing quality levels. Since software is on the critical path for many engineered products, and for internal business systems as well, the new approaches are starting to affect global competition and attract widespread international interest. It can be hypothesized that success in mastering software quality will be a key strategy for dominating global software markets in the 21st century.

Session A1: Software Management, 10:15-11:45 am, TTC Auditorium

•

t.

Т

ŧ

ł

)

}

Rodemy Ashby, Sandia National Laboratories

The Right Rock: Finding and Refining Customer Expectations

Figuring out what the customer wants, making sure the team understands the customer priorities, and negotiating what the customer can have for what they want to pay sets the scene for project success or failure. Getting a clear understanding of the political landscape (can't tell the players without a scorecard), and what is most important to them is essential. The people who will be using the system you produce, and those paying for it are rarely the same, and both must be satisfied for your project to be considered successful for the long term. Ways to bring internal differences of opinion to the fore, and flush out misunderstandings while educating the customers and project team about the cost of different decisions involves creating a vivid, shared understanding of how the target, completed system looks and operates. Approaches to these problems that I've found useful include 1)Erika Jones Organization Charting, 2)Customer Interviews, 3) Quality Functional Deployment and modifications with other "matrix-type" decision-making tools, 4)Creating an initial system acceptance test document, keyed to the requirements as requirements are negotiated, 5) Rapidly-Prototyping an example to show the customer, and modifying it per request if you have a configurable system and/or 5)Create the User Manual first. TH illustrate the methodology and tool use with project examples.

David Harris, Sandia National Laboratories

TCAMS Lessons Learned

The overall objective of the Technical Control, Automation, Maintenance, and Support (TCAMS) system software is to facilitate the operation of the communication center within the Commander in Chief (CINC) Mobile Alternate Headquarters (CMAH). The software consists of about one million lines of source code and draws heavily upon industry standards such as Ada, SQL, Unix, and X-Windows. Several technical decisions that were made during the design and implementation of TCAMS went awry. This presentation attempts to provide insight into the root causes for these wrong decisions with the hope that these insights can lead to a better understanding of the software development process. An overview of the TCAMS project including some measures of the software complexity is included as introductory information.

Joseph R. Schofield, Jr., CQA, Sandia National Laboratories

The Next Silver Bullet - Or Just Another Shot in the Foot?

Repeated promises of productivity and quality improvements have seldom materialized with the introduction of new technologies. Marginal incremental improvements in productivity have become accepted as the norm. Joe shares a model that explains the unintended outcomes of technology hopping as well as how to extend the investment in a technology. Further implications exist for maintaining and improving the ability to manage the software development process as measured with instruments such as the Capability Maturity Model. The notion of the "inflight magazine syndrome" only exacerbates efforts to stabilize and maximize our use of technology. This work was recently published as the lead article in Managing System Development.

Session B1: Software Testing, 10:15-11:45 am, Bidg 822 Rooms A&B

Debra Sparkman, Los Alamos National Laboratory

A Working Testing Process

Argus is an automated security system deployed at 4 DOE and DoD facilities across the United States. Argus is composed of 3 major subsystems including over 20 software and firmware products. This paper describes the processes performed for testing the Argus Security System. The primary focus is on the independent testing activities. A brief description of unit, integration, and system testing performed by the development staff will be presented. Independent testing is conducted by the Quality Assurance team using a separate test system. The independent testing process is a practical approach to implementing independent testing for an existing software-based system undergoing major enhancement development. The primary focus of testing is based upon system level regression testing, major feature enhancements and new product testing. Test planning is conducted prior to each testing activity. This planning is based upon risks associated with the degree of modifications and their impact on the customer operational systems. The testing process is completed by the generation of a test report summarizing the testing activities. This work was performed under the US Department of Energy by Lawrence Livermore National Laboratory under Contract No. W-7405-Eng-48.

Nancy A. Storch, Lawrence Livermore National Laboratory

Testing the Design and Operations of a New Badging System

In response to a DOE mandated order to rebadge the Laboratory, efforts got underway to modify, replace, or adapt three major hardware and software systems. On a prior project, it had been helpful to conceptualize a complex system by gathering all interested parties together and systematically walking thorough a paper process description. However for the rehading project we needed to do more than conceptualize the end system. We needed to test operational aspects and integration of the systems with users in an environment similar to the actual deployment environment. This became a full-scale mock exercise of rebadging. Each system was in a different state of development. One was somewhat operational and in testing, one had a working prototype, another was in the lowfi paper prototype stage. Also, they were being developed by different teams which rarely interacted with each other. These teams were focused on designing, implementing and unit testing within their system. Therefore, traditional integration and system testing of the combined systems was still a long way off. We wanted to save development time through early identification of issues, integration and operational problems, as well as usability problems. In the mock exercise we had 22 participants, who came from the development teams, operations and maintenance, user groups, managers and customers. Observers were selected both from within and outside the project. Observation posts were identified to include coverage of both individual system operation and overall operations. Operational scenarios based on prior rebadging experiences were developed with hypothetical person's to be rehadged. Realistic artifacts were acquired or created. Message and data communication between systems was modeled using paper messages and records. Logistics were handled to turn a mothballed badge office into the futuristic badge office of the exercise. The exercise took place over three half days. By the third day, we had created a variation on the operational scenarios which held promise for a more streamlined operation. We also gained insights on the interactions and communications between the systems and a list of important issues, problems and action items was produced. This talk will focus on our approach to testing and discuss its costs and benefits within the software development life cycle.

Dwayne Knirk, Sandia National Laboratories

Establishing a Three-Way Agreement: Specification, Code, Test

After we complete software testing, what do we know and what don't we know about the subject computing system? What kinds of system tests will further reduce our ignorance about the suitability and correctness of the computing system for its application? Software-intensive systems are expected to work in a particular environment to bring about desired effects in that environment. To accomplish these effects, the computing system must have a variety of interactions with that environment. Its capabilities and features are directed to establishing a variety of relationships between those interactions, including stimulus-response, constraint, and historical reference. To establish such relationships are the services provided by the computing system. The given environment and required effects in the problem are collectively documented as Problem Requirements. The computing system interactions and services are documented Behavior Specification. The relationship between these two sets of

information is an explicit and verifiable *behavior design* task. The Behavior Specification characterizes a computing system independently of its application context. It provides a single reference point for all decisions of software architecture and implementation as well as for test case and testware architecture and implementation. Had we error-free development and testing processes, we should expect specific behavioral equivalencies between the pairs (specification, code) and (specification, test). To the extent these processes are not perfect, we may have defects in our code, our tests, or both.

This presentation explains the logical implications of the behavioral equivalencies, and interprets them in operational terms. It described how testing provides a means of comparing software and testware behaviors and evaluating their behavioral equivalence to the source specification. An integrated testing approach is devised for identifying deviations from the desired equivalence. The approach provides specific guidance for test design, test execution, code design, instrumentation and data collection, and evaluation of test results. The presentation concludes with a summary of what can be known through this logic-based testing approach and what remains to be examined in final system testing. The ultimate goal is validating the behavior of the resulting system through measuring its effects in the application environment.

Session C1: Software Quality for Scientific Applications, 10:15-11:45 am, Bidg 822 Room C

John Ambrosiano and Robert Webster, Los Alamos National Laboratory

Software Quality and Process Improvement in Scientific Simulation Codes

Today the reliance on high quality software is so important that standards for quality assurance are an integral part of software development in both the public and private sectors. Yet as a community, research scientists have not entirely embraced these methodologies and indeed are often leery of them. Is the problem with scientists, or with the standards? As the quest for excellence in software is extended to government research activities, we must understand this phenomenon and either modify how SQA standards are introduced to the scientific community, or understand why they are inappropriate, and if inappropriate, how to modify them. A salient aspect of research software development is that it usually involves a high degree of novelty and risk in the beginning. Only later, after evolving through a series of prototypes, are concepts considered sound enough to be turned into production software. This sometimes leaves scientists at a loss in deciding when to introduce their products into the SQA process. Too early and progress toward developing useful new concepts is impeded. Too late and high quality may be impossible to assure. In this paper we apply process analysis and knowledge acquisition methods to study the evolution of simulation models for nuclear technology applications from seminal prototypes to production design codes. Using use-case scenarios and interviews, we will build a model of the simulation software production process. We will also try to understand how the expert judgments of the scientists involved contribute to their ranking of a software product's quality and readiness for production. We will compare the results of this analysis to the practices recommended to attain SEI's CMM level 2 certification. In doing so we will try to answer the following questions: Which of these software development activities best fit a SQA model such as the SEI CMM and which do not? Is there a modification of the CMM that allows research scientists to more easily introduce their software at some appropriate stage into a standard SQA methodology?

÷

Ł

Edward W. Russell, Lawrence Livermore National Laboratory

The SQA of Finite Element Method Codes used for Analyses of Pit Storage/Transport Packages

This presentation will describe the implementation of the SQA requirements of DOE/AL, Quality Criteria (QC-1), Revision 8, July 1995, for Finite Element Method (FEM) codes used at the Lawrence Livermore National Laboratory (LLNL) for conducting design and confirmatory analyses on pit storage/transport package designs. This work satisfies the requirements of the Defense Technologies Engineering Division (DTED) Quality Assurance Policy and Plan for software management of activities associated with high risk, commensurate with the LLNL risk-based graded approach of SQA implementation. Element 14.0, "Software Quality Assurance," of QC-1 dictates the following requirements: (1) organization, tasks, and responsibilities; (2) verification and validation; (3) configuration management; (4) software documentation; and (5) reviews and audits. The FEM codes controlled by this program are utilized for structural and thermal analyses. As an example, DYNA3D which was originally developed at LLNL in the late 1970's, is a nonlinear, explicit, three-dimensional FEM solid and structural mechanics code for analyzing transient dynamic responses. Element formulations include onedimensional truss and beam elements, two-dimensional quadrilateral and triangular shell elements, and three-

б

dimensional continuum elements. Many material models are available to represent a wide range of material behavior. Sophisticated contact interface capabilities are available, such as frictional sliding and single surface contact. The size of DYNA3D is roughly 100,000 lines of code with 700 subroutines.

The SQA implementation for FEM codes is guided by the commercial standard, *ISO 9000-3: Guideline for Application of ISO 9001 to the Development, Supply and Maintenance of Software*, with increased SQA formality as necessary to satisfy the requirements of the nuclear standard, QC-1. The IEEE SQA standards and guides were consulted for guidance on format of the SQA Plan and associated specifications. The IEEE recommendations were tailored for this application to meet the requirements of the governing document, QC-1. The requirements within the DTED QA system to maintain and control high-quality software include the following documentation for FEM codes: SQA Plan, Requirements Specification, Design Description, Configuration Management System (CMS), and Verification and Validation Report. The CMS uniquely identifies and controls code versions and changes, as well as all pertinent baselines, procedures and documentation. Validation is accomplished by using a suite of analytically and experimentally validated benchmark problems.

Orval hart, Los Alamos National Laboratory

Software Quality Assurance at the Weapons Engineering Tritium Facility

The Weapons Engineering Tritium Facility (WETF) at the Los Alamos National Laboratory began construction in. 1982 and finally received authorization to go on-line in 1991. It was the first nuclear facility to receive authorization under Admiral Watkin's increased formality-of-operations. Due to the many changes in DOE orders for nuclear facilities, the facility took longer than would be expected to get on-line. First it was "yes, we'll grandfather you in under the old regulations", then it was "no, you will have to meet the new regulations". The WEIF went through several Readiness Assessments (then called Safety Appraisals) and the Operation Readiness Review before finally receiving approval to start operation. The WETF is unique, in that it was the first nuclear facility to place what was previously administrative procedures (interlocks, etc.) into software that was monitoring and controlling major operational aspects of the facility. The Instrumentation and Control System is designed to be inherently safe, i.e., if any of the computers controlling the facility fails, the systems will fail safe. That is, all valves are closed, all pumps stopped, etc. The facility cannot be operated in this mode, but is left in a safe state. Backup procedures allow for the safe restarting of the facility. Many of the operational systems are automatic in their nature, i.e., the ICS takes immediate action when an 'operational' abnormality occurs. Operation of the facility, in general, is performed from Operator Consoles in the Control Area, as opposed to through switches or hands-on in glove boxes. Due to this new method of operation, where software is involved in almost all operation and surveillance of the facility, the DOE was 'extremely' apprehensive about how all this was to work. This presentation will discuss the Quality Assurance program that was adopted to assure that the WETF could be operated in a safe and reliable manner.

Session A2: Software Engineering Processes, 01:30-03:00 pm, TTC Auditorium

Michael Bell, Lockheed Martin Energy Systems

Function Point Count Adjustment by Means of Scaling Touched Function Points

The talk presents an adjustment method to function point analysis that will quantify the work effort involved in a software enhancement project in terms of function points. The technique allows direct comparison of the magnitude of work with the magnitude of functionality change, which is also measured and expressed in terms of function points. The adjustment method is based on effort data that are ordinarily readily available, avoiding complex and costly data collection requirements or subjective judgments. The technique accounts for software development activities that are not directly measured by function point analysis. The adjustment may be used with attribute analysis to predict and baseline a wide range of software development efforts.

Stewart Meyer, Westinghouse Savanuah River Co.

Using An Automated Code Management System To Improve Configuration Control Practices

Using a configuration management tool (software library) is not something new, several organizations and Sites use them. There are monerous tools commercially available, some claiming to be extensible and easy to customize. We took a very simple tool and added a front end to it. This front end is the interface to the software libraries and

shields the users from knowing the command language of the tool. In addition, the front end enforces the configuration control policies as set forth in the QA plans and procedures. The methods are then consistent across organizations and software products that are managed using this system as a tool. The front end is a developed product that may be used in other areas at the Savannah River Site, or other Sites, assuming the base system components are available. Although this system is used by one section at SRS, it could be available for use by others, without further investment in hardware. The key processes to improve were:

1

۲

ı

ł

I

ł

ł.

1. Identification of baselines; 2. Methods for verification of patches in a process control; environment; 3. Performing concurrent development in a controlled environment; 4. Methods for implementing periodic verification; 5. Configuration audits.

Outline of this presentation:

1. Description of deficiencies in previous software CM methods; 2. Description of methods and practices changed to foster improvements; 3. Description of SCMS system architecture and software tools; 4. Functional description of the SCMS from a user perspective relative to CM practices.; 5. Discussion on how key processes were improved.

Karen Jefferson, Terry Porter & Todd West, Sandia National Laboratories California

Software Engineering and Graphical Programming Languages

In a Work for Others project for the Air Force, The Advanced Atmospheric Research Equipment (AARE) software team used National Instruments' LabVIEW (a data flow graphical programming language) to control hardware used to collect samples of airborne particulate and gaseous species. Along with developing control and data collection software, the customer required MIL-STD-498 processes and documentation. This talk will discuss the processes and tools developed to support this project from the requirements to testing phase. In addition, unique aspects of the processes specifically tailored to graphical programming languages (such as coding standards, coding documentation, and configuration management) will be presented.

Session B2: Internet WEB Applications, 01:30-03:00 pm, Bidg 822 Rooms A&B

Kevin Hill, Pantex Plant :a Kie

Internet Strategies for Engineersure

The tools available on the Internet have the potential to help engineers reduce costs and increase productivity. As the amount of information available increases, so does congestion. Thus the Internet may be a victim of its own popularity. Strategies for effective use become necessary. How can an increase rather than a decrease in productivity be achieved? A survey of engineers' Internet usage is the first step in the search for ways to optimize time on the Internet. Two methods are used to advance this search. The first is the interpretation of survey results and follow-up questions. The second is via literature review. Standard search methods in conjunction-gifth human networking can make the Internet a more productive tool. Concerns which have restricted Internet-usage, such as reliability of sources, and unwanted leaking of information are addressed. Survey results and analysis-provide a forum to initiate a discussion of this powerful tool's (the Internet's) impact on engineering efficiency and software quality.

David Leong & Fran Current, Sandia National Laboratories

Exploiting the Intranet: A New Architecture for Enterprise Information

The Intranet is an architecture for viewing information within the enterprise. This architecture is based upon the World Wide Web standards. With the global Internet as a proving ground, this architecture is proving to be a very formidable information system for corporate uses. One of the strongest features of an Intranet is its inherest cross platform support. Applications are functional on PCs, Macintosh, and UNIX platforms. The basic purpose of most betranets today is the electronic delivery of corporate documents. These documents are typically of a static nature; corporate policy, manuals, newsletters. With the presentation capabilities of a web browser, compelling documents with integrated text, graphics, sound, and even video can be delivered via the Intranet. Hypertext links allow documents to be integrated in a way that makes knowledge even more accessible when compared to print media. Database access through a web interface is also a very powerful tool to the corporation. Query access to MIS systems typically living on the mainframe can now be made available to everyone on the Intranet adopting a three tiered client-server strategy, the web can become a graphical interface to legacy systems. Now the corporation's electronic phone book, human resource information, and financial reports can be delivered via a web

browser. Creating interactive web interfaces involves additional technologies. Security, workflow, and the 'Javas' (JavaScript from Netscape, and Java from Sun). In the area of security, authentication and anthorization are very integral to client-server applications that allow the user to update information. Transactional based workflow is also necessary to route task requests among workgroups in the enterprise. Standard HTML forms offer a stateless user interface. By using Java and JavaScript, one can create applications that establish connections and provide field level event handling on the presentation tier of the application.

This new paradigm for delivering information is not without its share of challenges. Cultural and political barriers exist that must be addressed with the same vigor as the technical challenges. An enterprise solution must have input from users within that enterprise. It is necessary to show the users how the enterprise Intranet can make their daily job easier. The enterprise web (Intranet) is a scalable productivity tool for the corporation that will enhance the way employees do their job.

Jeanie Negin, Sandia National Laboratories

"Rightsizing" Software Quality for a Web Services Organization This presentation describes variations of software engineering and project management as applies to an organization that is supplying services for Sandia National Laboratories' Intranet on a cost recovery basis.

Session A3: Software Process Improvement I, 03:15-04:45 pm, TTC Auditorium

Don Schilling, AS/FM&T

Quest for Excellence 1996: Reaching for the Stars

In the Spring of 1995, a need for software process improvement arose when DOE requested that certain software be handled as product. A solution was needed quickly to meet critical production schedules. This presentation summaries the actions and the processes that were followed in developing and implementing a solution for FM&T to handle product software. It discusses the Total Quality improvement process used and the outputs which were developed. The presentation is based upon the presentation given at AlliedSignal in the Quest for Excellence competition. The Quest for Excellence is a corporate-wide competition designed to show case process improvement. The team won the Teamwork Award for their efforts in defining a system which worked successfully and minimally impacted critical production schedules. This presentation also ties in with the tutorial of how the Nuclear Weapons Complex projects should be handling software as a product in response to Engineering Procedure EP401099. It shows one sites struggle in defining a workable process to meet customer expectations.

Don Rathban, AS/FM&T

Command Media System at the Kansas City Plant (KCP)

The Kansas City Plant was certified to the ISO9001 Standard in April 1995, following a successful audit by Third Party Auditor, Det Norske Veritas (DNV). The KCP has also successfully passed three six-month periodic audits by DNV subsequent to receiving certification in 1995. A new on-line Command Media System was developed and implemented to help ensure control of the documents associated with the KCP business processes. This control is demanded by the International Organization for Standardization to receive ISO9001 certification. The new on-line system is based upon the KCP Business Model. New Process Descriptions (PDs) and Work Instructions (Wis) were created by the KCP Process Owners for each process and released in the Command Media System. The development of the KCP Business Model and the new Command Media System will be discussed during the presentation, including how to access the system and structure of documents within the system. Also to be discussed are the operational structure in place to manage Command Media and proposed improvements to the system in 1997.

Michael Tiemann, Headquarters Department of Energy

Departmental Information Architecture

The Information Technology Management Reform Act of 1996 requires agency Chief Information Officers (CiO) to develop, maintain and facilitate the implementation of sound and integrated information technology architectures. Notwithstanding this act's formalization of this recent requirement, the Department of Energy's Designated IRM official, the Assistant Secretary for Information Management, decided well over a year ago to

establish a Departmental or enterprise-wide Information Architecture. As described in the published document the Department of Energy Information Architecture, Volume One, The Foundations, dated March 1995, the Departmental Information Architecture is a high level, principles and standards based framework within which the majority of programmatic, organizational and field site architectures should be developed and implemented. It is intended to be a template that can guide all information management acquisitions, activities, projects, developments, solutions and implementations. In order to help achieve this goal additional documents have been written to further explain and define the Architecture. Two additional volumes, Baseline Analysis and Guidance, (Information Architecture Volumes Two and Three, respectively) have been published to describe the current or defacto Departmental Information Architecture and to provide specific guidance on the establishment of Information Architectures within other organizational components of DOE. The intent is that they will be treated as nested organizational subarchitectures within the overarching Departmental Architecture. The Baseline Analysis document identifies many of the challenges facing the Department in regard to the divergent, often incompatable, obsolete, or non interoperable technologies and systems currently deployed as well as the duplication and redundancies, inheritant in the applications and data structures. The Guidance document provides useful guidelines for architectural activities in all life cycle phases for DOE and its partners and stakeholders. In addition, there are several architectural standards related documents being published and widely distributed. Presently there are numerous architectural efforts underway at various sites and within several of the major programs. It is the intent of the Office of the CIO to support these activities and to grow this approach further throughout the entire DOE community.

This presentation will summarize the above documents and related actions and activities to date regarding the Departmental Information Architecture Program and explain the future directions as the Departmental Information Architecture becomes the Chief Information Officer's central component in the comprehensive Departmental Information Management Strategy.

Seasion B3: High Integrity / Formal Methods I, 03:15-04:45 pm, Bidg 822 Rooms A&B

Larry J. Dalton & Marie-Elena Kidd, Sandia National Laboratories

Meeting the High Integrity Software Needs of Today and Tomorrow

Quantifiable measures of the reliability, safety and security for software-based systems remains an elusive goal even after decades of research. Such systems continue to be a major source of safety and security catastrophes. These catastrophes include the of loss of life, environmental or economic damage, and loss of public confidence. In spite of these catastrophes, the usage and complexity of software-based systems in high-consequence applications is continuing to increase. This growth, with the associated safety and security risks, presents a national challenge to the R&D community. Sandia National Laboratories established a High Integrity Software research project in 1995, to begin to address the challenge. The first of two research areas, the Correctness Track, is focused on creating the 'ability to create software that is "correct by construction." Research projects include advanced concepts for the capture of software specification/requirements, validation through intuitive and visual reasoning and mathematics for correctness preserving transformations covering all steps from specifications to executable code. The second research area, Systems Immunology, is directed towards in-situ techniques and technologies to enable real-time fault detection and safing control (fault response). Systems Immunology research projects include Software Event Execution Reliability (SEER), Digital Isolation and Incompatibility, and Top-Down Fault Analysis of Microprocessor Systems.

Victor Winter, Sandia National Laboratories

An Overview of the AST Software Construction Methodology

AST is a formal method that is being developed within the High-Integrity Software (HIS) project at Sandia National Laboratories. AST stands for Abstraction, Synthesis, and Transformation. Within AST, abstraction, deductive synthesis, and transformation techniques are used to enable the automation of a significant portion of the software construction and verification process. Furthermore, within AST the impact of human involvement is limited to such an extent that it can be formally verified. In AST, the role of synthesis is to construct abstract algorithmic solutions to problems from nonalgorithmic specifications (e.g., precondition and postcondition pairs).

2

This is accomplished by using a sophisticated search engine such as an automated reasoning system to resolve (or remove) the nondeterministic choices that are present in the initial nonalgorithmic specification. Complementing

Presentation Abstracts: Wednesday, April 2 1997

synthesis within our methodology, the role of refinement transformation is (1) to optimize solutions that are obtained in the synthesis step, and (2) to introduce low-level (e.g., machine oriented) algorithmic details for the purpose of (ultimately) producing a machine executable implementation satisfying the original nonalgorithmic specification. Currently, AST is restricted to a somewhat well-behaved subset of reactive systems that we refer to as single-agent reactive systems. Because the burden placed on the synthesis portion of our methodology can be enormous we have found it useful to distribute the synthesis process over an abstraction hierarchy. In order for this approach to succeed, the abstraction hierarchy must have the property that a solution at one level of abstraction "benefits" or "can be used to guide the construction of" a solution at the next lower level of abstraction within the hierarchy. In essence, what is going on is that an algorithmic skeleton is being synthesized at one level of abstraction and is then in some sense "passed down" to the next level in the abstraction hierarchy. This process continues until a machine executable algorithm has been obtained. An undesirable consequence of this approach is that the synthesized algorithms tend to be sequential in nature (i.e., completely parallel or concurrent solutions cannot be readily synthesized in this framework). Fortunately, it is well within the capability of refinement transformations to take a sequential specification of a problem and to then transform it into an efficient parallel solution. This talk gives an introductory overview of AST as well as a brief example of how transformation techniques can be used to compliment synthesis.

Alex Yakhnis & Vladimir Yakhnis, Pioneer Technologies

Towards Automated Construction of Dependable Software/Hardware Systems

Many observers have recognized that software/hardware systems built by Government and by Industry can be very complex. It may be difficult to establish dependability and functionality of such systems. Here are some of the questions that existence of such systems raises. (1) How a software/hardware system should be documented in order to be understood by users and customers of various backgrounds? (2) What should be established in order to conclude that the system is acceptable? (3) Finally, since the system intent is often evolving in the course of system design and use, how should we modify the system to reflect this evolution while preserving the system dependability? Here are some of the approaches which are presently used in Industry in order to resolve the above questions: (1) Presenting a system as a hierarchy of models where the levels of the hierarchy would represent various levels of abstraction. Then an observer could look only at the levels of hierarchy that do not have details that are of no interest for the observer. Another approach to document a system is the object oriented approach. Here, systems are understood through understanding of individual objects from which the system is composed and of interactions among objects. Usually, the approaches are not combined. Also, thus far applications of objectoriented approach were mostly limited to the software-only system components. (2) Exhaustive testing that system behaviors satisfy the requirements. The problem here is that exhaustive testing is not possible even for moderately complex systems. An approach to overcome this is to formalize system requirements, to accurately model the system that is being constructed, and to produce a mathematical proof that the system model satisfies the requirements. However, so far, this was done with respect to system components only. Moreover, correctness proofs are usually not applied to several software constructs, e.g. communication among objects. (3) Maintaining system requirements, models, design, and simulation information in a single data base capable of containing many system versions. However, such a data base alone would not insure that the next version would be as dependable as the previous one. In this talk we will describe a direction of work on how to get better answers to the above questions on the basis of mathematical modeling, formal methods, and multi-agent strategic approach. These methods are aimed to achieve industrial strength automation of system specification, design, correctness proofs, and maintenance without exhaustive testing. Mathematical modeling and formal methods are beginning to be recognized in Industry as promising approaches to deal with high complexity of systems. The formal methods groups have been formed at Intel, Motorola, and HP.

Session A4: Software Process Improvement II, 08:30-10:00 am, TTC Anditorium

L

;

. . .

1

ł

1

i

Cathy Kuha, AS/FM&T

AlliedSignal Capability Maturity Model Assessment & Improvement Processes

This presentation provides a summary of the processes used by AlliedSignal to assess progress against the Software Engineering Institutes Capability Maturity Model and the use of this assessment data to plan and implement organizational process improvements. AlliedSignal corporate has committed to achieve CMM Level 3 at sixteen of its key business units within the next three years. This strategy is a key component in an effort to develop a competitive advantage in the aerospace business. What's unique about this initiative is that it is being applied to Information Systems. Staff at the AlliedSignal Aerospace Center for Process Improvement and the AlliedSignal Corporate Information Systems group have developed the methods and materials to assist business units in this strategy. Six certified SEI examiners have been trained to conduct progress assessments and supporting material have been developed. Included in this material is a process guide for using assessment results to plan and drive organizational improvement. Each business unit is scheduled for a formal assessment every 6 - 8 months. Quarterly self-assessment metrics are provided by each business unit and are used to track progress. The presentation focuses on the continuous improvement cycle implemented at the Kansas City site as a result of repeated assessments and planning.

.

Ann Stewart, Lockheed Martin Energy Systems

Lessons Learned on Utilizing the SEL/CMM in the Federal Government Work for Others Environment

Data Systems Research and Development (DSRD), a division of Lockheed Martin Energy Systems, Inc., has developed a specific approach in applying the Software Engineering Institute's Capability Maturity Model (SEI/CMM) that has been successful in our customer focused environment of research and development within the federal government. This approach is based on establishing an orderly and understood infrastructure consisting of three major building blocks, controls, processes, and information. This infrastructure is sustained through a strong quality program emphasizing technical, peer, and management reviews and quality audits and surveillances. This paper describes the tactical application of this approach and DSRD's experiences and lessons learned in three years of implementation.

Gail Benefield, Lockheed Martin Energy Systems

"SWIM" Your Way to Software Quality

A company quality improvement effort has many aspects. At Lockheed Martin Energy Systems at Oak Ridge, a software development methodology called Software WorkPackage Methods (SWM) has been created and can be considered part of the company's quality improvement efforts. SWM is a methodology for managing, developing, and supporting information system projects and applications. It is composed of methodology guidelines, role definitions and assignments, and work packages. The work packages are in the form of work breakdown structures suitable for project estimating, planning, and management. SWM provides development and support processes which are customizable, yet repeatable. It keeps pace with new software development methods and techniques and provides automation support for the project estimating, planning, and management.

Session B4: High Integrity / Formal Methods II, 08:30-10:00 am, Bidg 822 Rooms A&B

Mikhail Auguston, New Mexico State University

Debugging Automation Tools Based on Event Grammars and Computations over Traces.

Dynamic program analysis is one of the least understood activities in software development. A major problem is still the inability to express the mismatch between the expected and the observed behavior of the program on the level of abstraction maintained by the user. We propose to design software testing and debugging automation tools based on assertion language concepts as well as on precise program execution models. We are developing a PARFORMAN language for the description of computations over execution histories of target programs that provides a basis for tool development for assertion checking, debugging queries, execution profiles, and performance measurements. We use assertion language mechanisms, including event patterns and aggregate operations over event traces, to describe typical bogs and debugging rules, and to evaluate debugging queries. An event grammar provides a sound basis for assertion language implementation via target program automatic instrumentation. These tools and methods may be useful for software testing, debugging, documentation, and

maintenance of software systems. Our approach is nondestructive, since assertion texts are separated from the target program source code and can be maintained independently. Assertions can capture the essential dynamic properties of a particular target program and can formalize the general knowledge of typical bugs and debugging strategies. Event grammars may be designed for sequential as well as for parallel programs. Examples of assertions and debugging rules for run-time detection of bugs and bug localization are presented. We have developed a prototype implementation of the assertion checker and debugging rule evaluator.

Marie-Elena Kidd, Sandia National Laboratories

A Method for Critical Software Event Execution Reliability in High Integrity Software

When high consequence systems rely on software for critical control functions, they require high integrity software. A major concern of high integrity software is ensuring the faithful execution of critical software driven event execution sequences. To meet system performance criteria, high integrity software must execute correctly and reliably. In addition, in the presence of transient hardware or software faults in both normal and abnormal environments, safety and security objectives must be maintained. A reliable, repeatable method and application techniques are needed to address these issues. Our technical approach involves an in-situ (embedded in the software) dynamic (run-time) fault detection and mitigation method for ensuring critical event execution sequences in high integrity software. Our method is based on deriving a mathematical description of the critical software controlled event execution sequence from a software model or the software requirements, embedding check points and update points based on that mathematical description into the target code, and adding a software module that implements the functionality of the underlying mathematical model. This extra software is added to the target code to verify that the correct software event execution sequence is maintained.

John Sharp, Sandia National Laboratories

Business Rule Enforcement Via Natural Language Modeling

The topic of my presentation will be business rule enforcement using Natural Language Modeling. A well defined procedure will be explained that allows subject matter experts to specify requirements and then be held accountable for them. I will convey a fundamental truth: That requirements can always come in the form of precisely analyzed, elementary natural language sentences.' Requirements include both facts that result in tables for populating data and business rules that do not change the table structure, but they do restrict the population of otherwise good facts in existing tables. A brief review of analysis results will now be discussed to allow you to understand a portion of the capabilities of this procedure. The following sentences all require external data to populate the instances of knowledge that is desired to be maintained.

Professor has degree in subject.

Course requires minimum degree level in subject

Professor teaches course.

Referential integrity applies, in that populations of the third sentence must be from known populations of professor and course in the first two. These sentences cannot enforce the business rule that a professor must be allowed to teach a course before he can be assigned to teach the course. I define this requirement as a "business rule" because no other fields are needed to store the data than appears in the previous three, but the rule can be enforced by starting with the derived sentence:

Professor is allowed to teach course.

This sentence is a derived fact (an SQL query can be established with appropriate triggers) and a set theory rule can be applied to restrict the population of the third sentence. This rule is:

The professor teaching a course must be a subset of the professors who are allowed to teach that course.

All "business rules" can be written as either direct set theory constraints against facts that are externally populated or as derived fact(s) and set theory constraints against other facts or derived facts. The benefit of Natural Language Modeling is that all of the experts and users can understand and be held accountable for the specification of the design because it always exists as a set of understandable sentences. Transformations of this knowledge set can be made into any graphical technique (including relational and object-oriented methods) but I do not know of any graphical presentation that can handle all of the knowledge captured.

÷...

Session A5: Software Process Improvement II, 10:15-11:45 am, TTC Auditorium

Larry Desonier, Sandia National Laboratories

Guns for Hire - Experiences of Quality Software Development Under the Gun

In today's software development environment, a major concern is the quality of the software. Sometimes getting the quality boxes checked seems to take precedence over implementation and delivery. There exists a way to both perform rapid development and have a quality product. There is a saying that 80% of the work gets done in 20% of the time, and the rest may never get finished. The question here is simply can quality software be developed when (1) 80% of the dollars are spent, (2) only 20% of the work is complete, (3) there is 6 weeks to delivery, and (4) no code has yet been written (and the team estimate is many months to code completion). This is just the situation for a "Guns for Hire" team. In some organizations this would be known as a type of "Skunk Works" or software "Swat Team." Our experience has shown that with the right size team, the right skills mix of individuals, and some disciplined development practices, quality software can be developed and projects can be saved. This discussion will reflect on projects accomplished in just this manner: projects developing user interface or command console software, a PC-based graphics display for alarm annunciation, material and personnel tracking systems, a taxi-way monitoring system, and others. This would not be possible without an experienced team, standard development practices, actually reusing code (yes, it is possible), and strictly disciplined development practices. The successes of this process paradigm is why the "Guns for Hire" team is continuously in demand.

Bruce Johnston, Panter Plant

The Year 2000 Challenge: A Project Management Perspective

Today we are faced with the biggest threat to computing ever discovered. As the year 1999 makes its final tick into the year 2000, many time-sensitive business applications like accounting, payroll, project management and many, many more will either completely fail or make disastrous mistakes. Why will this happen? In the 1970's and early 1980's when data processing shops were buying mainfiame computers by the truckloads, the high cost of memory persuaded programmers to drop the century digits from a date field to save two bytes of memory. Although shortsighted, this practice was universally accepted because these early computer applications were not expected to be in operation today. Using only two digits for the year 1996, for example, is represented simply "96." This means when the year 2000 arrives, tens of thousands of old software programs still in use will think the year is 1900. If the doornsday predictions hold true only half of the worlds computer applications will be completely fixed or replaced before December 31, 1999. This will be a real challenge: finding, changing, and testing date parameter software changes and the challenge will be an even greater Software Quality Assurance problem for legacy programs. This paper will address the year 2000 challenge from a project management perspective and give insight into managing the project of the century.

Curt Holmes, Lockheed Martin Energy Systems

Year 2000 Awareness

The Date 2000 challenge has been referred to as both a technical problem and a business risk. It has also been called the single largest information technology project which corporations and government agencies will undertake in the next several years. Current estimates for the cost of remediating Date 2000 software problems in the U.S. range between \$600 billion to \$1000 billion, and are increasing. The problem will affect all hardware platforms and all software systems in various ways and with unpredictable results. On average, organizations are finding that over \$0% of their existing applications portfolio is impacted by two-digit year date processing (i.e. 19xx). Some systems will shut down, while others will corrupt data and generate spurious output. In all cases, the business operational risks, resulting from the failure of internal operating systems, far out weigh the potential cost of remediation. The purpose of this presentation is to create an awareness of Year 2000 issues, promote collaboration among DOE sites, and propose electronic sharing of resources to save money in infrastructure and software resources costs.

ι

Session B5: High Integrity / Formal Methods II, 10:15-11:45 am, Bldg 822 Rooms A&B

John Hare, AWE UK

ISO and Software Quality Assurance

Emerging International Standards now promise a global approach to Software Quality Assurance; ISO/IEC 12207 provides a framework for Software life cycle processes that has already attracted the attention of both US and UK customers. The ISO 'SPICE' standards give international weight to the concept of self-assessment, and a model that could take the SEI CMM world-wide. Previously our customers have independently developed their own standards, which include QC-1, AQAP 150 and DetStan 05-95. Whilst ISO9000-3 can be adopted for assessment, this is non-mandatory and has not been well received in the US although widely used in Europe. TickIT, the scheme for third party assessment, could refocus on ISO/IEC 12207. This presentation reviews customer requirements and the new International Software Standards, with particular emphasis on ISO/IEC 12207 and SPICE. It is concluded that ISO Standards will become a dominate driver for Software Engineering, and could now succeed in promoting a world-wide approach.

Larry Rodin, Pantex Plant

Licensing and Certification of Software Professionals

This report presents information on software engineering certification programs, licensing of software engineers, reasons to become certified, certification as a condition of employment, the body of knowledge and examination structures for the certification programs, and an overview of the Institute of Electronic and Electrical Engineers recommendations for software engineering as a profession.

The Software Quality Assurance Subcommittee of the Nuclear Weapons Complex Quality Managers completed a Work Item to research software-related certification and licensing efforts and provided status reports to the Quality Managers. A white paper was a significant result of that work item and this presentation has been updated to reflect changes in the licensing and certification processes.

Certification is a voluntary process administered by a professional society. Licensing is a mandatory process administered by government. Two professional organization have been identified as having or developing certification programs, and one state has developed legislation for a licensing program:

- The Institute for Certification of Computer Professionals (ICCP) has two levels of certification Associate Computing Professional, and the Certified Computing Professional;
- The American Society for Quality Control has implemented its program for Certified Software Quality Engineer;
- New Jersey is the only state identified as actually enacting software development legislation, their licensing program covers "software designers".

Included in the presentation are considerations and implications for licensing and certification. What problems are we solving by having licensing and certification. Equal Employment Opportunity (EEO) laws will be discussed to address issues such as: can certification testing being considered discriminatory; or can certification as a condition of employment be considered discriminatory.

Michael Lackner, AS/FM&T

Operational Excellence (Six Sigma) Philosophy: Application to Software Quality Assurance

The Kansas City Plant, as part of AlliedSignal Aerospace, has committed fifteen individuals to each receive four months of training in Six Sigma and at least a year in the position established as a Blackbelt. Six Sigma is a philosophy of doing business encompassing the methodologies of defect prevention (versus defect detection) through the use of statistical tools, i.e., process mapping, design of experiments, and process controls. Business includes providing any product or service. Continuous improvement to the way business is performed is achieved through the identification of optimal target values in products and processes, and the reduction of variation around those targets. An overview of the tools and training will be discussed, along with the application to the processes included in Software Quality Assurance.



Session Z: Keynote Tutorial

Dr. Dave Parnas NSERC/Bell Industrial Research Chair in Software Engineering McMaster University Ontario Canada

Session	Title
Z 0	"Design Through Documentation: The Path to Software Quality"
21	"Inspection of Critical Software"
Z2	"Exercise & Discussion"



Keynote Tutorial

Professor David Lorge Parnas, PhD holds the NSERC/Bell Industrial Research Chair in the Communications Research Laboratory, Department of Electrical and Computer Engineering at McMaster University in Hamilton, Ontario, Canada. His primary area of interest is to promote the discipline and body of knowledge to Software Engineers as practiced by engineers in other fields. By studying the problems of software engineering since 1965, Dr. Parnas has developed principles and methods that have value to real world problems. In recognition of his accomplishments, he has received numerous honors, including election as a Fellow of the Royal Society of Canada and a Fellow of the Association for Computing Machinery.

Dr. Parnas will share his experience and knowledge by leading the three workshop/tutorials described on the next page.

David Lorge Parnas NSERC/Bell Industrial Research Chair in Software Engineering Communications Research Laboratory Department of Electrical and Computer Engineering McMaster University Hamilton, Ontario Canada L8S 4K.1 Phone 905-648-5772 FAX 905-648-5943 Email parnas@qusunt.crl.McMaster.CA

Z0: April 1 1997, 09:00 - 11:00 am, TTC Auditorium "Design Through Documentation: The Path to Software Quality"

In traditional engineering design, a series of documents precedes the actual construction of the product. These documents permit review and analysis, then after revision, serve as input to the next phase. When the (inevitable) errors are discovered and changes are required, the design documents already on file are updated and reviewed again. Each new refinement is reviewed against the previous documents.

In software design this "waterfall" method is almost never applied. Although it is appealing, practitioners are not able or willing to write precise documents. Instead, they write vague blurbs that are useless to those charged with the next steps and cannot be subject to rigorous analysis.

We will describe how precise, complete, and testable documents can be produced for software and the ways that these documents can contribute to an improved software process.

Z1: April 1 1997, 01:00 - 03:00 pm, Bidg 823, Breezeway "Inspection of Critical Software"

Software is devilishly hard to inspect. Serious errors can hide in a software product for years. People are hesitant to employ software in safety-critical applications. Many companies are finding correcting and improving software to be an increasingly burdensome cost.

This talk describes a procedure for inspecting software that consistently finds subtle errors in "mature" software, software that is believed to be correct. The procedure is based on three key ideas:

- The software reviewers are active not passive
- Reviewers focus on small sections of code.
- Reviewers proceed systematically so that no case and no section of the program gets overlooked.

During the procedure, the inspectors produce and review mathematical documentation. The mathematics allows them to check for complete coverage; the notation allows them to proceed systematically and in small steps.

Z2: April 1 1997, 03:00 - 035:00 pm, Bldg 823, Breezeway "Exercise & Discussion"

Participants will be given a small program and will apply the documentation and inspection methods to them. This will be followed by a discussion of previous experiences in question and answer format.

Design through Documentation: The Path to Software Quality

David Lorge Parnes

NSERCEst) Industrial Research Chair in Software Engineering Communications Research Laboraiery Department of Electrical and Computer Engineering Michineer University Franktion, Ontwile Consta LSS 451

Abstract

In traditional engineering design, a series of documents precedes the actual construction of the product. These documents permit review and analysis, then after revision, serve as input to the next phase. When the (inevitable) errors are discovered and changes are required, the design documents already on file are aplated and reviewed again. Each new refinement is reviewed against the previous documents.

In software design this "waterfall" method is almost never applied. Although it is appealing, practitioners are not able or willing to write precise documents. Instead, they write vague, burbs that are useless to those charged with the next steps and cannot be subject to rigorous analysis.

We will describe how precise, complete, and testable documents can be produced for software and the ways that these documents can contribute to an improved software process.

> Communications Research Laboratory Software Engineering Research Group "connecting decay with practice"

MoMaster University a

Why is Software So Often a Problem?

Developers consistently underestimate the difficulty

systematically, identify and record requirements.

explicitly design, document and review software

These steps are standard practice for all engineering

Communications Research Laboratory Selements Englworing Research Ortop "connecting Beety with Station"

hold reviews of the requirements document,

· carefully inspect all designs and programs.

The steps are not taken for software because,

of building software for long-term use.

They write software rather than design it.

deligner: aller

They do not:

spucture,

products other than software.

"The code is self-documenting!"
"Software is just a set of instructions."
"Anyone who knows the language can program."

"Software is easy!"

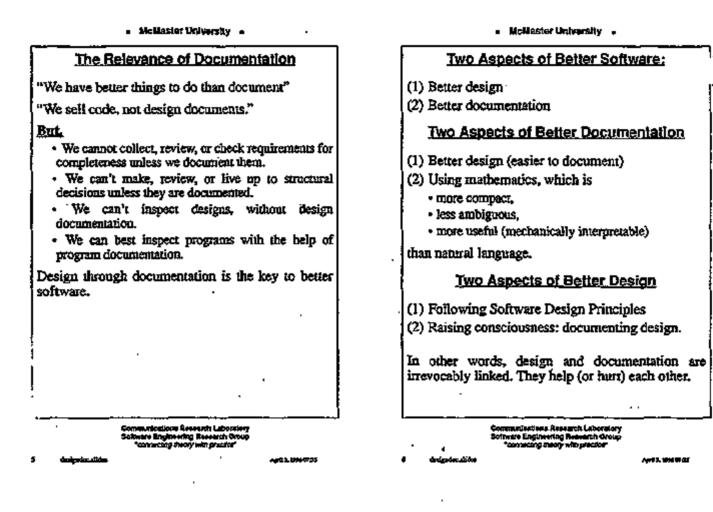
Famous last words!

4493,29660905

McMaster University =

The Goal: Better Software at Lower Cost Software is a collection of software components, Nobody can build products as one big "blob" Everyone wants to re-use software components. "Components are junk!" (industry leader) What's the problem? Components are hard to re-use (hidden assumptions). Components have complex interfaces Components are not well documented. The design process does not emphasize these issues. material and Research Laboratory Science Engineering Research Gost "connecting heavy with procior" Active protected Mollaster University Why Don't People Apply Engineering Discipline to Software? Some don't have an engineering education. (2) Some don't think it's necessary. (3) Some don't know how to do it. Why don't we demand that software people have ? radiacillary standard Experience shows that it is necessary. In this talk I want to focus on how to do it. Communications Resistent Laboratory Software Englaneting Research Occup features in any with practice

4 designation



Writing Down Requirements

The most costly errors are those made early in the process - they are the hardest to change.

Misunderstandings about requirements lead to early mistakes.

Progutaness need to be told what is needed.

They must also be told what is subject to change.

Requirements must be subject to review.

Safety reviews of software must be based on a previously agreed statement of requirements.

Maintenance actions must be based on requirements.

None of these things is possible unless we have a <u>written</u> statement to work with.

That <u>written</u> statement must be precise and complete.

- McMaster University -

What's Wrong with Requirements Methods?

We think of requirements as a set of elements, each element being one requirement.

Consider three such requirements.

- · The output must be an integer.
- The output must be positive.
- . The output must not be zero.

Consider an alternative formulation:

· The output must be a natural number

These are equivalent - one requirement or three?

We cannot count requirements or list them?

If we try, we have no hope of checking for completeness, consistency, correctness.

There is a better way, based on the basic model used in control theory.

Communications Research Laboratory Sufference Englanding Research Group "exercised Restry with practica" Communications Retearch Laboratory Satistary Englanding Receptor Carup Connecting Receyority practics

412.05000

McMaster Onlyershy =

How to document system requirements?

The first step is to:

Identify monitored variables $(m_1, m_2, \bullet \bullet \bullet, m_p)$.

Identify controlled variables (c1, c2, ***, cp).

The primary monitored variables are things <u>outside</u> the system whose values should influence the output of the system. Examples:

- customer meter reading.
- steam temperature
- -time of day

The primary controlled variables are things <u>outside</u> the system whose values should be determined by the system. Examples:

- what the operator sees
- what appears on a bili
- control positions

This is only the beginning, but for many projects you cannot even find a complete list of these variables and there is no agreement on what they are.

Communications Research Laboratory Balances Engineering Research Group "contacting Restynetic practice"

-43, gen 62,51

مانانىمىدراما ا

McMaster University

Bringing Time into the Picture

All of these variables can vary with time.

For each scalar variable, x, denote the time-function describing its value by " x^{t_0} .

The value of x at time t is denoted "x^t (t)".

The vector of time-functions $(v_1^t, v_2^t, ..., v_n^t)$ will be denoted by " v^{t*} .

Contrary to the statements of some computer scientists, there is no problem dealing with "real" time.

McMaster University • _

Monitored and Controlled Variables Will Be Added During The Design Process.

It is inevitable that the need for additional variables will be discovered as we get into detailed work.

Further, <u>new</u> monitored and control variables are created during the design process.

The <u>primary</u> monitored and controlled variables are outside the system.

Sometimes we want to monitor the system itself, i.e. measure things that did not exist before the system was built.

Sometimes we may even want to control (adjust) parts of the system.

As the design is developed, we may add these monitored and controlled variables to the requirements document.

It is essential that the document be updated. Otherwise reviewers and maintainers are lost.

> Companiestions Research Laboratory Solowers Engineering Research Croup "connecting Receipting Society"

in inighteration

4912 1999 0745

NeMaster Univoteity

Bringing Math Into our Tool Kit

The implementors need to know the following relations:

Relation NAT:

domain contains values of m¹, sange contains values of s¹,

(m⁰, g⁰) is in NAT if and only if nature permits that behaviour.

This tell us what we need to know about the environment.

Relation REO:

- domain contains values of m¹, range contains values of g¹.
- (\mathbf{n}^i, g^i) is in REQ if and only if system should permit that behaviour.

This tells us how the new system is intended to further <u>restrict</u> what NAT(ure) allows to happen.

If we can describe these relations, we have our system requirements written down.

We can get the "scary" math out of the documents by using the right notation.

Communications Research Laboratory Software Engineering Research Group "connecting Patry with Stacket" Communications Research Laboratory Software Engineering Research Group "converses metry with pressor"

فحيتمه 1

Why Use This Approach?

- For all the "motherhood" reasons that we try to find the requirements first.
- (2) Because we can check for completeness.
- (3) Because we can check for consistency.
- (4) Because we have a precise description.
- (5) Because we have a reviewable document.
- (6) Because we can often simulate the system.
- (7) Because the design can be based on the document.
- (8) Because the programming goes much faster.
- (9) Because the programmers work consistently and do not duplicate each other's work.
- (10)Because we will discover ways to simplify the system.
- (11)Because we can build monitors for testing or supervising the system.

Why not?

(12)Because it requires some training.

(13)Because it is a risky front-end investment that slows down the initial part of development.

Converting for the Converting Sector of Cooperating Sector of Coop

13 independente di terreta de la constante de

April 2 Distances

McMaster University

How can we document system design?

 i^{t} denotes the vector valued time function $(i^{t}_{1}, i^{t}_{2}, \cdots, i^{t}_{r})$ with one element for each of the input registers

 \underline{o}^{t} denotes the vector valued time function $(o_{1}^{t}, o_{2}^{t}, \bullet \bullet \bullet, o_{q}^{t})$ with one element for each of the output registers

Document the following relations Relation IN:

domain contains values of m^t, range contains values of j^t

• $(\mathbf{m}^i, \mathbf{i}^i)$ is in IN if and only if input device permits that behaviour

It must be the case that $domain(IN) \supseteq domain(NAT)$

Relation OUT

- domain contains the possible values of <u>o</u>^t
- + range contains the possible values of $\underline{c}^{\rm L}$
- (o¹, c²) is in OUT if and only if output device permits that behaviour

Communications Research Laboratory Solitative Englishering Research Comp "consecuto favory with proton"

N datigadat dilas

April 2, 1046 (PAS

1 3 100 CONT

MeMaster University

When Can We Skip System Design?

Sometimes the I/O devices are simple and we can have simple relationships between the controlled and output variables as well as between the monitored and controlled variables.

In that case, we can use the systems requirements document as a software requirements document.

Many applications have this property.

In some, we can cheat and mix the two.

McMaster University

Dividing the Software to Conquer Complexity

Small modules are easier to understand, if the interfaces to other modules are simple.

To keep interfaces simple, "hide" the details inside the module.

Use the requirements documents to help structure the software:

- + Some modules hide the requirements (REQ)
- some modules hide software decisions (which are not in the requirements document).
- Some modules hide the hardware (IN, OUT) These modules are support software.

These modules "create" virtual:

- data structures,
- devices,
- "actors",

"objects" that do part of the job.

It is at this stage that we have the best chances for reuse - but we must document the interfaces.

Computing one Research Laboratory Software Engine wing Research Group Communications Research Laboratory Software Engineering Research Group "connecting Recey with practice" MeMaster Doiversity

Documenting Module/Object Interfaces (1)

It is wise to design software by designing a set of objects.

- · Each object is implemented by a module (a set of programs) using a data structure that is "hidden from" (never used directly by) programs outside the module.
- Changing the state of the object, or getting information about the object's state, is only done by invocations of programs from the module.
- An object is a finite state machine.
- The input alphabet of an object is the set of operations one can perform upon an object.
- The output alphabet of the object is the set of values that can be returned by such operations.

The state of an object can be hidden.

Describing or specifying objects is very different from describing or specifying programs.

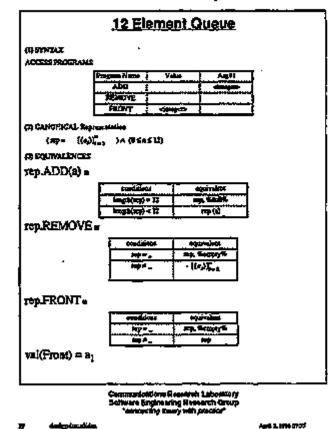
Hiding the state means that we must discuss event sequences, but it makes future changes easier.

Conventional Research Laboratory Software Engineering Research Group "connecting theory with practice"

Aug 1, 200 (80)

LT designation a links

MoMaster University



McMaster University •

Documenting Module/Object Interfaces (2)

Black-box interface descriptions must be written in terms of (input, output) sequences (traces).

- A <u>trace</u> of a finite state machine is a finite; sequence of pairs, each containing a member of the input alphabet and a member of the output; alphabet.
- A trace, T, is considered possible for machine M, if M could react to the sequence of inputs in T by emitting the sequence of outputs in T.

Descriptions and specifications of objects can both be written as predicates on classes of traces.

These predicates are the characteristic predicate of an extension function/relation.

We organise our descriptions in terms of:

A canonical abstract state representation, and

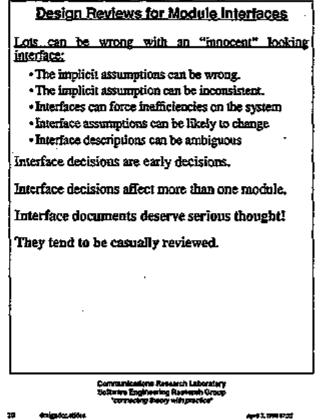
single event extensions of those traces.

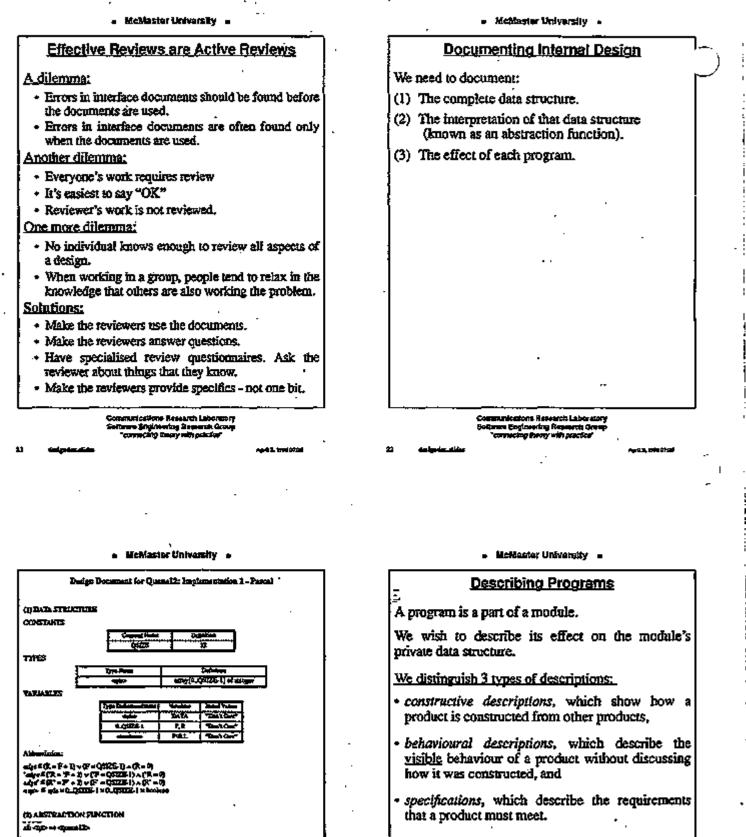
Result: a systematic, reviewable reference document

ontractications Research Laboratory attante Raylmenting Greatest Great "consecuty theory was precipit

Ang 3, 1994 (1995)

McMaster University





In my view this is a very important distinction that is ignored by the "formal methods" community.

Company feelborn Research Laborates Software Engine using Research Orou "consorcing Resay with practics"

entrop (1474) - 70- ... , entrop

10.00

Constantiantical Research Laborator Software Engineering Research Grou "constrainty theory with processor"

or (linteration)

(-++++*F0113.4#1470)

بالكتار ومشغ

- ALAY PERSON AND A DESCRIPTION OF A DES

April 1, 1990 (7.31

Relational Program Descriptions and Specifications

Users need to know the relation between the starting values of variables and the final values of variables.

Users need to know the starting states for which the program is guaranteed to terminate.

We base our work on Harfan Mills' ("Cleantoom") program function, but

- Represent the function in a more readable tabular format.
- · Deal properly with non-determinism.
- Carefully distinguish between relations as specifications and relations as descriptions.

It is possible to produce short, readable specifications of programs and review them before writing the actual code.

This forces designers to think about issues that they tend to overlook (such as error response).

Communications Revisedh Laboratory Salpune Englavoring Ressanh Droup "connering duray asli pischaf

April 3, (2004) (2)

Molfaster University

The "Laws" of Programs

Do Software Engineers have laws for programs that correspond to Kirchoff's laws for circuits?

Yes!

The basic laws of programs are essentially the axioms of the algebra of (LD-)relations.

If you, accept the fact that LD-relations provide adequate descriptions of program behaviour, sequential execution is composition.

The laws are the classic results about relations.

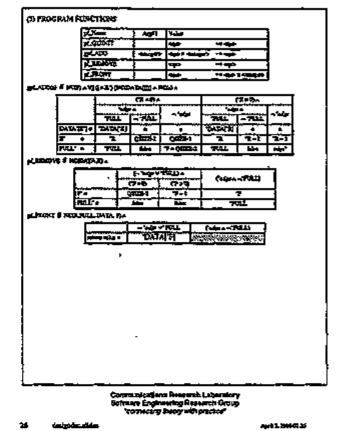
These laws allow you to find behavioural descriptions of constructed programs if given:

the constructive description of those programs and,

the behavioural descriptions of the primitive programs.

With these laws, all reasonable specifications and descriptions are compositional. Composition is not Conjunction.

McMaster University



MeMaster University

Imperfection of Documents?

When engineers work with physical products they <u>must</u> use imperfect implementations of abstract specifications.

With software, imperfection is not always necessary but it may be convenient and acceptable.

The imperfections must be "bounded" and explicitly limited in their applicability.

For example, we may ignore the limits on representations of numbers because we only work with a limited range of numbers.

It is important to include this in the specification.

No new mathematics is needed for this. Implication does the job.

The use of mathematics in engineering does not imply a belief in perfection of programs or maths.

April 2,200 (2012)

What New Notation do we Need?

Although the mathematics is old, and the abstract notation for defining things is old, the applications are new.

We have to describe relations and functions that have non-heterogeneous ranges and domains and can have a discontinuity at arbitrary points.

We have found a variety of tabular notations to be useful.

Ryszard Janicki, has found new ways to unite these tabular notations.

Jeff Zucker and our students are implementing tools for transformations.

We are trying to:

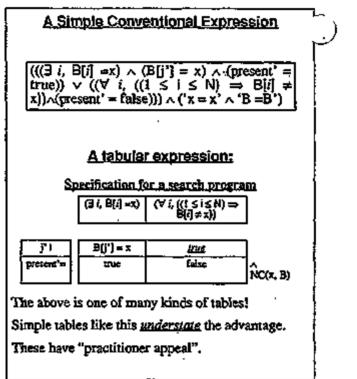
- Make the documentation easier to produce
- Make the documentation more useful

Communications Research Laboratory Settemes Engineering Research Group "connecting Sectory with gracifor"

29 designéer-Méri

April 3, 1996 (#195

McMaster University



Communications Research Laboratory Sufframe Engineering Research Casup "survivoing 2000/problem"

McMaster University

Inspecting Programs

Its the code that "hits the road."

Getting the requirements right, the structure right, the interfaces right, etc. are all important but we have to check the code.

The same review principles apply.

- · Make the reviewers use the documents.
- Make the reviewers answer questions.
- Have specialized review questionnaires. Ask the reviewer about things that they know.
- Make the reviewers provide specifics not one bit.

We want to compare the completed programs with previously reviewed specifications.

We ask the reviewers to produce descriptions.

We then show that the descriptions match the specifications.

It's hard work but it produces results.

- · We get good documentation for future use
- We find errors in the best industrial code programs that were considered correct.

Sonraceferiers: Research Laborator Informite Englanding Remarch Group "columnary Junity with propints"

MoMaster University +

Is it Teachable/Learnable/Practical?

Its the way to start - first year engineering students have learned to read and implement from specs.

Tabular notation - no theoretical advantage, but a great practical advantage.

Short courses introduced these ideas to the nuclear industry in Canada. They now teach their own.

People can apply the inspection technique after a 3 - 4 day coarse.

Critical Mass in a company is essential. Writers without readers are useless.

There is lots of room for improvement. We will identify these faster if you work with us.

> Communications Research Unbornion Seitures Engineering Research Group "connecting theory with precies"

3) designates allega

22 ánimta á

61.WHERE

-03, 29440-34

Sets for Describing Programs

Everything about digital computers can be explained in terms of finite sets; the set concept is viewed by many as the most basic concept in mathematics.

A set is a collection of elements from a previously defined set (sometimes called the universe).

The elements in the universe must be known before other sets are defined. Every application of set theory must begin with a careful description of the Universe from which it's elements are drawn.

Sets drawn from different universes cannot be compared.

Set elements are assumed to have previously defined attributes.

The famous anomalies can be avoided.

Call it dull set theory.

Communitations Research Laboratory Software Englaceding Research Drawp
"and being Dairy with practica"

فتجمع والمتحجمة

2/19

Petersoy 34, 391

	McMaster	University	
--	----------	------------	--

	Notation for scient
{ x ,y,z}	enumeration -a set containing x,y, z
l suc	zh that
{x l <cond satisfies th</cond 	ition>)The set of elements such that x are condition.
A⊆B	A is a subset of B (could be identical)
$A \subset B$	A is a subset of B and smaller than B
AυB	set of elements in either A or B
AnB	set of elements in both A and B
A – B	set of elements of A that are not in B
(B) (the comp	set of elements in Universe not in B lement of B)
X∈A	X is an element of A
{} an	empty set
Only com	bine sets from the same Universe.
Even em Universe.	pty sets must have an associated
	Communications Research Laboratory

Communications Research Laboratory Settent's Englaneting Research Group "saturatory recoy with product" 2/19

kçeçetili is

Newsy K. M

McMaster University

Relations

What is a relation $(e,g, \ge, <, =)$? A set of ordered pairs.

What is the domain of a relation? The set of elements that appear as the first element of a pair in the relation.

What is the range of a relation? The set of elements that appear as the second element of a pair in the relation.

One need not enumerate all the pairs to describe a relation!

If R is a relation and $(x,y) \in \mathbb{R}$, we can write $x \mathbb{R} y$.

McMaster University

Examples of relations

Both elements taken from the set of real numbers.

- (1) $A = \{(x,y) | x > y\}$
- (2) $B = \{(x,y) | x = y\}$
- (3) C = A ∪ B
- (4) $D = \{(x,y) \mid x \times x = 4\}$
- (5) $E = \{(x,y) \mid x \neq y = 4\}$

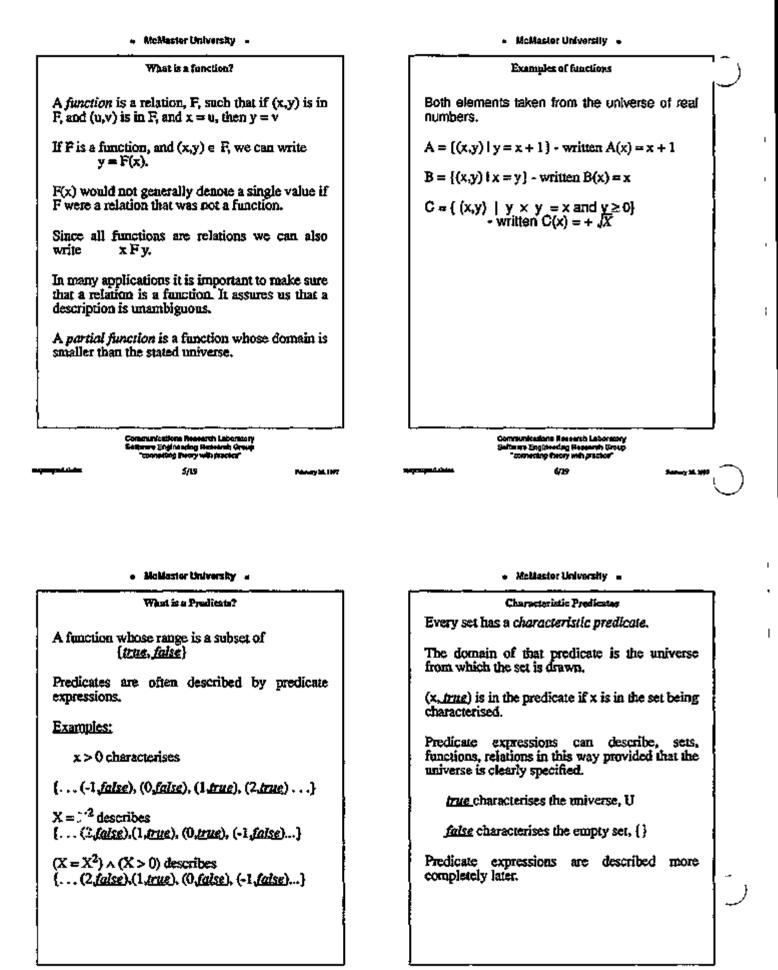
Communications Research Laboratory Software Engineering Research Group Community Encody web practice 4/19

nunicatione Research Labora

re Endra alión Responds C

in proposition of the

W II. 1997



Communications Research Laborancy Settemin Englaneting Research Group "concepting many with practice" 7/29

er 16. 1997

Communications Research Laboratory Safatain Singlineering Research Group "Connecting Recey to practice" 8/19

Characteristic Predicates Describing Relations

 $\{(x,y) \mid x < y\}$ described by x < y

$\{(x,x') | x'='x+1\}$ described by x' = 'x+1

The use of predicate expressions in this way requires clearly stated conventions about the universe and the naming of the elements of an ordered pair.

'x can be read "x before" or "x left".

x' can be read "x after" or "x right".

A predicate expression is not a predicate.

A predicate expression is not a set.

A predicate expression is not a function or relation, Predicate expressions can describe;

predicates sets

> functions relations

> > Consequences in the second statements of the second second

فالد فحجات

9/19

Boliaster University

Definition of Predicate Expressions

Built-in functions and predicates are named:

To simplify the presentation we shall assume that all functions and relations have simple nemes.

 $f_1, ..., f_k$ are the names of functions (sets)

 $R_1, ..., R_m$ are the names of the characteristic predicates of relations.

•	McMaster University
---	---------------------

Summerv A relation is a set of pairs (2-tuples). The set of values that appear as the first element of a pair is called the *domain* of that relation. The set of values that appear as the second element of a pair is called the range of that relation. A function is a relation such that for any given element, x, in its domain, there is only one pair (x,y) in the function. If (a,b) is in the function F. "F(a)" means b. often called "the value of F at a", may include niples, the may make sense to write "F((a,b))". "F((a,b,c))", and "F(F((a,b,c)))". -Functions whose domain is smaller than the universe are called partial functions Most of the functions that arise in software development will be partial functions. A predicate is a function whose range contains no members other than true and false. •For any set, X, the characteristic predicate of X is a predicate whose domain is the universe from which X is drawn, and whose value, for b, is true

> munistions Research Laboratory tana Engineering Research Group "anneting Resay with procees" 10/19

if and only if b is a member of X.

er 14, 1993

 McMaster University **Definition of Predicate Expressions** 2 Terms are constructed from: A finite set of mathematical variables, $x_1, ..., x_n$ A finite set of constants, C The constants are strings. Each constant represents one member of the universe, U. "V" stands for a comma separated list of terms (see below). A function application is a string of the form f;(V). A term is either a constant, a variable, or a function application.

diere Research Laboratory are Engineering Research Droi connecting Reary with practice"

n 144 diana Sara teate Engl wing Research Ge connecting many with press 12/19

1039

Definition of Predicate Expressions

A primitive expression is a string of the form R_j(V).

Nothing else is a primitive expression.

All of our expressions will be built of primitive expressions.

Note that primitive expressions, since they denote predicates, will always evaluate to either true or false.

> manufactions Research Laboratem nare Englassing Research G Connecting Resty with practice

McMaster University

The Meaning of Predicate Expressions

(1) If t is a constant representing t' (a member of U), the value

of the term 1 for assignment a, (written "val(t,a)"), is t', (2) I is a variable, x_k, the val(La) is the value specified for that

(3) if t is a function application, f₄(V), we must evaluate each of

V* denotes the result of this evaluation

We distinguish the following three cases:

the terms in V until we have obtained the values that they

(3a) if V' is in the domain of f_{b} , val(t,a) is $f_{b}(V')$, (3b) if V' is not in the domain of fip val(t,a) is <u>not</u>

(3c) if any of the elements of V' is not defined. the value of the function application is <u>not defined</u>.

Evaluating terms:

specific assignment.

varisble in a.

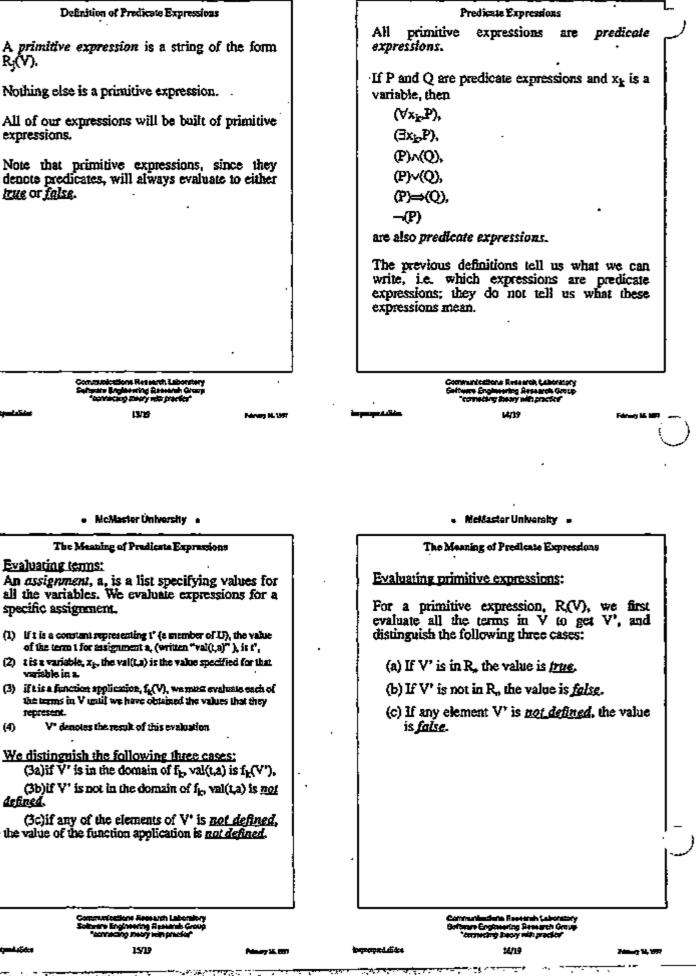
represent.

(4)

defined.

13/39

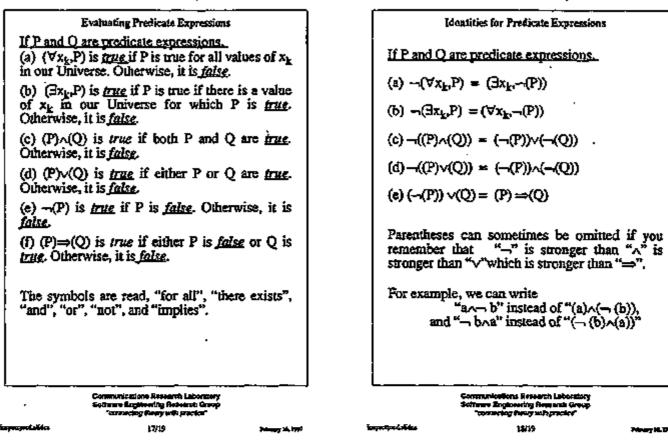
MeMaster University



uniostions Seesarch Labor incoding Research In Decity with price ere Engin

15/19

NoMester University



McMaster University +

- M. Den

MeMaster University

Examples of Predicate Expressions $((x>0) \land (y=\sqrt{x})) \lor ((x\leq 0) \land (y=\sqrt{-x})) (1)$ $((x>0) \Rightarrow (y = \sqrt{x})) \land ((x \le 0) \Rightarrow (y = \sqrt{x})) (2)$ $((y = \sqrt{x}) \vee (y = \sqrt{-x}))$ (3) $(\exists i, ((1 \le i \le n) \land (A[i] = X)))$ (4) $(\exists i, ((1 \leq i \leq n) \Rightarrow (A[i] = X)))$ (5) $((1 \le n) \land (\forall i, ((1 \le i < n) \implies (A[i] \le A[i+1])))$ (6) Exercise Try to write English statements corresponding to the above. - A-أنسنة طورته ngineering deserved cong freesy with pract Getoware Engine a. a.

20715

and a list of

17 M. HW

Software Inspections We Can Trust

David Lorge Parnas

NSERC/Bell Industrial Research Chair to Software Engineering Communications Research Laboratory Department of Sincipiest and Computer Engineering Net/Inter University Hamilton, Ontario Canada LAS 4821

Software is devilishly hard to inspect. Serious errors can hide for years. Consequently, many are hesizant to employ software in safety-critical applications and all companies are finding correcting and improving software to be an increasingly burdensome cost.

This talk describes a procedure for inspecting software that consistently finds subtle errors in software, software that is believed to be correct. The procedure is based on four key ideas:

- All software reviewers are actively using the code.
- Reviewers exploit the hierarchical structure of the code rather than proceeding sequentially through the code.
- Reviewers focus on small sections of code, producing precise summaries that are used when inspecting other such sections.
- Reviewers proceed systematically so that no case, and no section of the program, gets overloaded.

During the procedure, the inspectors produce and review mathematical documentation. The mathematics allows them to check for complete coverage; the notation allows the work to proceed in small systematic steps.

> Communications Research Laboratory Schware Engineering Research Group "comparing decay with practice"

insperie allerer

MeMaster University

Why is Software so often a Problem?

Developers <u>consistently</u> underestimate the difficulty of building software for long-term use.

They write software rather than design it.

They do not:

- · systematically, identify and record requirements,
- hold reviews of the requirements document,
- explicitly design, document and review software structure,
- carefully inspect all designs and programs.

These steps are standard practice for all engineering products other than software.

The steps are not taken for software because,

- "Software is easy?"
- "The code is self-documenting!"
- "Soliwere is just a set of instructions."
- * "Anyone who knows the language can program."

Famous last words!

Communications Research Laboratory Solitente Engineering Research Group "connecting theory with practice"

i ingentialiter

Names y \$ 6, 2007 \$3.54

Férenz 16, 597 2004

McMaster University

Responsibilities of (Software) Engineers

- To thoroughly understand the properties of their products.
- To follow established rules of good practice when designing and building products.
- To apply accepted theory where it has been shown to lead to better, safer products.

Engineering is Not Management

The art of management is the ability to get things built without knowing exactly what they are.

The engineer is expected to thoroughly understand the properties of the product.

Software projects are hard to manage - especially if they are badly designed, but ...

Unless we have good Engineers, the best managers will not be able to successfully manage these products.

> Convententions Research Laboratory Saturno Segimeering Research Gravy "conversing theory with practice"

in a state of the state of the

Ferring 20, 1997 23 (4

MoMaster University

Why Don't People Apply Engineering Discipline to Software?

(1) Some don't have an engineering education.

(2) Some don't think it's necessary.

(3) Some don't know how to do it.

Why don't we demand that software people have appropriate qualifications?

Experience shows that it is necessary.

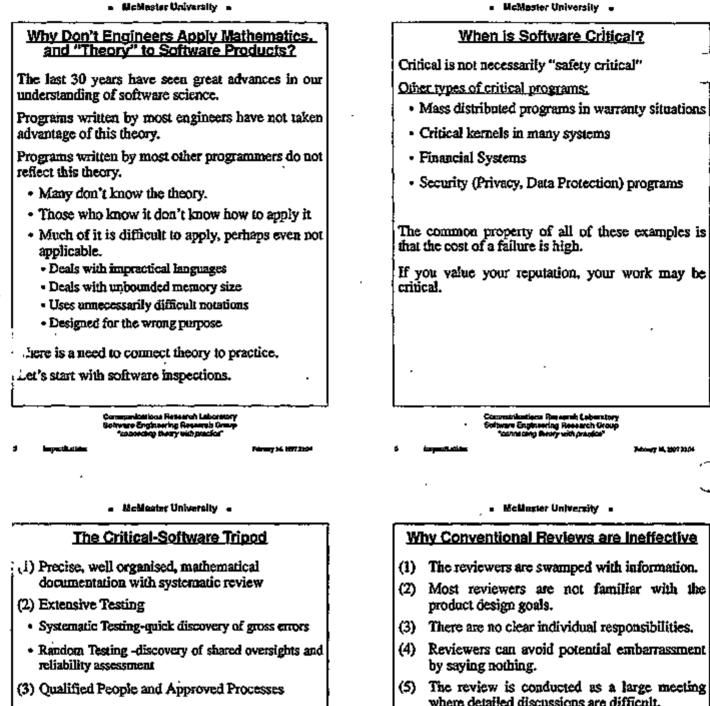
Why aren't software designers required to be Engineers?

Why do we continue to think of them as scientists and to educate them accordingly?

Conversional Research Laboratory Boltzare Englaneting Research Broup "bootecing theory with gradies"

interney 24, 2017 23,54

...



The Three Legs are complementary

The three legs are all needed.

The stool falls over if any leg is forgotten.

The third leg is the shortest.

his the shortest leg that we should worry about.

Today we discuss only leg (1).

zanieticny Reserve Laborate zwe Engineering Reserve Gro

14, 597234

(6) Presence of managers silences criticism.

(8) Specialists are asked general questions. Generalists are expected to know specifics.

review into a tutorial.

(7) Presence of uninformed reviewers may turn the

(10) The review procedure reviews code without

aus lections for such Laborrie was Englassifing Resource Gro

respect to structure. (n lines per hour)

(11) Unstated assumptions are not questioned.

Effective Reviews are Active Reviews

A dilemma:

- Errors in programs and design documents should be found before the documents/systems are used.
- Errors in programs and documents are usually found when the documents are used.

Another dilemma:

- · Everyone's work requires review!
- . It's easier to say "OK" than to find subtle errors!
- · Reviewer's approval is not reviewed.

One more dilemma:

- No individual can review all aspects of a design.
- When working in a group, people tend to relax in the knowledge that others are also working the problem.

Solutions:

- Make the reviewers use the documents.
- Make the reviewers document their analysis.
- Have specialised reviews. Ask the reviewer about things that they know.
- Make the reviewers provide specifics not just a bit.

Communications Research Laboratory Software Englacering Research Group "community theory with practice"

74mmy 16, 1977 21.04

eryan di katika

McMaster University

Parnas/NRL/AECB/AECL/Ontario Hydro

Focus on the engineering side.

Depend on hierarchical decomposition rather than sequential reading.

Use mathematical notations to provide precise descriptions rather than informal paraphrases.

Produce useful documentation as a side effect.

Proceed much more quickly if the documentation was already produced by the developers.

Insures that cases and variables are not overlooked.

Applies simple mathematics to check for completeness aspects.

Previous Work on Inspections

Best known approach Fagan - 1976.

Many followers - new book by Gilb.

Explicitly focus on the management aspects,

- Who should be there?
- · What are the roles of the participants?
- How long is a meeting?
- · How fast do you work?
- Forms for reporting errors?

Read the code in sequence and paraphrase.

Paraphrases are informal.

Most observers find these more effective than conventional reviews or walkfaroughs, but ...

... can we do better?

Companies Times Research Laboratory Software Englisheding Research Group

i ingena Datić

Fermy 20, 2077 20:54

• McMaster University •

Reviewing Design Documents

Base the review process on the nature of the document.

Begin by identifying desired properties,

Prepare questionnaires for the reviewers. Ask them questions that:

- make them use the document.
- make them demonstrate that the desired properties are present.
- ask for sources of information to support the answers to other questions.

For example:

- Ask reviewers to identify the domain of the program
- · Ask reviewers to identify "error" cases.
- Ask reviewers to explain why the behaviour required for each case is the desired behaviour.

For more information read [1].

Communications Research Laboratory Software Englaneting Research Group "connecting theory with practice"

February 14, 1997 \$250

Communications Research Laboratory Sotoware Engineering Research Comp

Teleny K, BR 204

11 ispailation

Inspecting Programs

It is the code that "hits the road".

Getting the requirements right, the structure right, the interfaces right, the documentation right, etc. are all important but we have to check the code.

The same review principles apply, viz:

- Make the reviewers use the material they review.
- Make the reviewers answer questions.
- Ask the reviewer about things that they know.
- Make the reviewers provide specifies.

We compare completed programs with previously reviewed specifications.

We ask the reviewers to produce precise descriptions.

We then show that the descriptions match the specifications.

It is hard work but it produces results,

- We get good documentation for future use.
- We find errors in the best industrial code programs that were considered correct.

Commun	deatloos Research Lai	borntory
	Engineering Research	
	tening theory with from	

logramik alid

Petrusy 34, 199733-34

McMaster University

Descriptions vs. Specifications

An actual description is a statement of some actual attributes of a product, or set of products.

A specification is a statement of all properties required of a moduct, or a set of products.

In the sequel, "description", without modifier, means 'actual description".

The following are implications of these definitions:

- A description may include attributes that are not required.
- A mecification may include attributes that a (faulty) τ , that does not possess.
- T₂₂ statement that a product satisfies a given specification may constitute a description,

The third fact results in much confusion. A useful definition has been lost.

> numbertions Research ners Engineering Re connecting Meany wit بغريا بلع Research Gr

McMaster University

Our Code Inspection Process

- (1) Prepare a precise specification of what the cod should do - a program function table.
- (2) Decompose the program into small parts appropriate for the "display approach" [2].
- (3) Produce specifications as required for the display approach.
- (4) Compare the "top level" display description with the requirement specification.

Observations:

- You can't inspect without precise requirements.
- Step 2 would already have been done if you use the display method for documentation.
- Step 3 is truly an active design review
- All reviewer work is itself reviewable.
- If you did not already have it, the by-product is thorough documentation.
- It's a bunch of small steps and very systematic.

Computationian Research Laboratory Software Engineering Research Group "moto-cite Meny withgradies"

(D. 41-4)

Many 14 107 3104

MoMaster University

Descriptions vs. Specifications

Any list of attributes may be interpreted as either a description or a specification.

Example:

"A volume of more than 1 cubic meter"

This could be either an observation about a specific box or, a statement of the requirements for a box that is about to be purchased.

A specification may offer a choice of attributes; a description describes the actual attributes, but need not describe the product completely.

Sometimes one may use one's knowledge of the world to guess whether a statement is a description or a specification.

Example:

"Milk, badly spoiled"

Guessing is not reliable. We need to label specifications and descriptions.

> Communications Research Lab Software Engineering Research "coloreday Meny with tract march Gro with the set

a 11, 227 2234

Do We Need New Semantics Theories For Programming 2

Not for the practical software engineering problems that I see.

I can find 30 year old theory that works for the problems that I will describe today.

Semantic theory has failed to describe real languages, but (in my opinion) the fault lies with the languages.

We do need improvements in:

- the notation used to describe actual programs
- the ability to describe behaviour in terms of the values of observable variables - nothing else.
- convenient ways to deal with all aspects of termination including non-deterministic nontermination.

What follows is mathematically equivalent to some <u>very old</u> ideas, but has some small practical advantages.

Communications Research Laboratory Salware Engineering Research Group San Meeting Advance with according

Ferrary 36, 829 242

7 60000-54

١

McMaster University

Using LD-Relations as Before/After Behavioural Descriptions (1)

Let P be a program, let S be a set of states, and let $L_P = (R_P, C_P)$ be an LD-relation on S such that (x,y) $\in R_P$ if and only if $\langle x, \dots, y \rangle$ is a possible terminating execution of P, and $x \in C_P$ if and only if P is guaranteed to terminate if it is started in state s. ¹ L_P is called the LD-relation of P

By convention, if C_P is not given, it is, (by default), Dom(R_P).

With this convention, our approach is upwards compatible with the "cleanroom" approach for dealing with deterministic programs.

¹ Please note that C_P is not the same as the precondition used in VDM [4]. S_P is the set of states in which the termination of P is certain.

Communications Research Laboratory Software Engineering Research Group "consecting Socry with Statics"

19 inspectilistics

. McMaster University ...

A Mathematical Interlude - LD-relations.

A binary relation R on a given set U is a set of ordered pairs with both elements from U, i.e. $R \subseteq U \times U$.

The set U is called the Universe of R.

The set of pairs R can be described by its characteristic predicate, R(p,q), i.e. $R = \{(p,q): U \times U \mid R(p,q)\}$.

The *domain* of R is denoted Dom(R) and is $\{p \mid \exists q \ [R(p,q)]\}$.

The range of R is denoted Range(R) and is $\{q \mid \exists p [R(p,q)]\}.$

Below, "relation" means "binary relation".

A limited-domain relation (LD-relation) on a set, U, is a pair, $L = (R_L, C_L)$ where:

 R_L , the relational component of L, is a relation on U, i.e. $R_L \subseteq U \times U$, and

 C_L , the competence set of L, is a subset of the domain of R_L , i.e. $C_L \equiv Dom(R_L)$.

Consumications Research Laboratory Soliware Engineering Research Croup "connecting theory with practice"

مة النيك بيون النيك بيون

Telony 18, 397 7134

Mollaster University

Using LD-Relations as Before/After Behavioural Descriptions (2)

The following follow from the definitions:

If P starts in x and x ∈ C_p, P always terminates; if
 (x, y) ∈ R_p, P may terminate in y.

• If P starts in x, and $x \in (Dom(R_P) - C_P)$, the termination of P is non-deterministic; in this case, if $(x, y) \in R_P$, when P is started in x, it may terminate in y or may not terminate.

 If P starts in x, and x ∉ Dom(R_p), then P will never terminate.

By these conventions we are able to provide complete before/after descriptions of any program but retain a simpler representation to use for those cases that arise most often.

> Communications Research Laborator Salarana Engineering Streamth Gran "concering Beory -isk practice"

and the second second

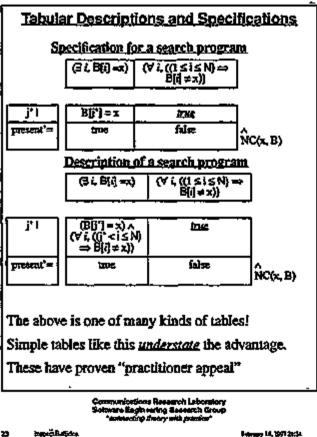
Theory 16, 1987 21:54

Petrony 14, 1977 24-26

Specifying Programs (1)

Settence Engin	ne Research Laboratory earlyg Research Group Sheary with principal Retrary 16, 20722-34
	element. If $S = \{L_S\}$ is a also call L_S a specification.
(2) P satisfies a specificati L _p satisfies at least or	
(1) P satisfies an LD-relation $C_S \subseteq C_P$ and R_P	
Let $L_p = (R_p \ c_p)$ to ine of Let S, called a <i>specificatio</i> LD-relations on the same $L_s = (R_s, C_s)$ be an eleme We say that	n, be a set of universe and
We can also use L1 specifications: Let $L_p \neq (R_P, C_P)$ be the d	D-relations as before/after
exhibited by a satisfacto	· }

McMaster University



here's Rations

Specifying Programs (2)

The following follow from the definitions:

- A program will satisfy it's own description as well as infinitely. many other LD-relations.
- An acceptable program must <u>not</u> terminate when started in states ouside Dom(RS).
- An acceptable program must terminate when stated in states in $C_{\rm S}$ $(C_S \subseteq Dom(R_P)).$
- An acceptable program may only terminate in states that are in Range(Rg).
- A deterministic program can satisfy a specification that would also be satisfied by a non-deterministic program.

Note. the following differences between the description and the specification of a program.

- · There is only one LD-relation describing a program, but that program will satisfy many distinct specifications described by different LD-relations.
- · An acceptable program need not exhibit all of the behaviours allowed by Rg (Rp G Rg).
- An acceptable program may be certain to ternainate in states outside $C_{S_{2}}(C_{S} \subseteq C_{P}).$

The intended use of each LD-relation (specification or description) must be stated explicitly!

CODE		Residen	ch Laboratory
	-	 uda a The	enerch Group
_			the second second

in the line of the

- 14, 1987 (1-14)

MoMester University

A Simple Example

(integer array H[1:N];

(integer c, integer a; $n \leftarrow 1$; # (n≤N→ (integer u; integer l; boolean p; $l \leftarrow 1$; $c \leftarrow 0$; # (u = l+n-l; (e≤N→(

```
(Integer i; i ← 0; p - true;
i/ (i < L(u - / + 1)+2 ) →
      (A[i+i] = A[u-i] \rightarrow (i \Leftrightarrow i+1; =)
£)
1
     (--p→ship (p→ c == o+1); / == /+1; = )
|u > N \rightarrow \oplus \rangle
Ë)
      \mathrm{H}[n] \Leftarrow q n \Leftarrow n + l; = )
```

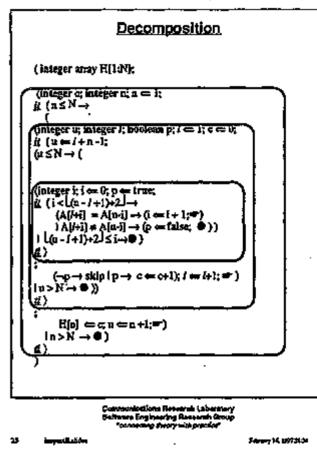
Constructions Research Laborat Software Engineering Research Grad "connected theory with practice"

#)

la>N→●)

NY 14, 1992 73-24

TEX



Displays: An Explanation

The top part of each display is the specification for the program in the middle.

The program in the middle is kept small by removing sections, creating a display for them, and including their specification in the bottom part.

The bottom part contains a specification of these invoked programs.

To check a display determine the description of the program in the middle, and see if it satisfies the specification at the top. In doing this, use the specifications of the invoked programs, not their text.

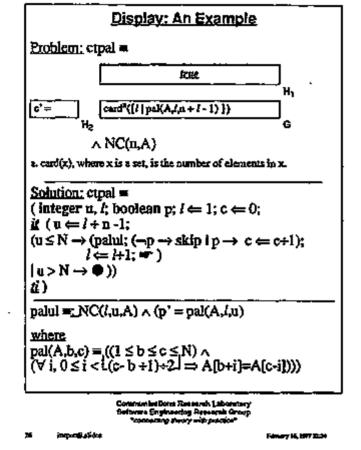
To check a set of displays, make sure that every specification at the bottom of one display is at the top of another. The exceptions:

- · standard programs
- primitive programs

Completeness can be checked mechanically.

Constructionalisms Research Laboratory Software Engineering Research Oroup "connecting shoory with practice"

McMaster University



UcMeater University

Structure and Inspection

Well-structured programs are easier to decompose. They can be decomposed by purely syntactic means.

Well-structured programs are much easier to inspect.

Inspection encourages good structuring.

Inspection suggests structural improvements.

Inspected programs are easier to maintain.

Modified programs need not be completely reinspected.

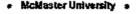
The cost of future maintenance is greatly reduced.

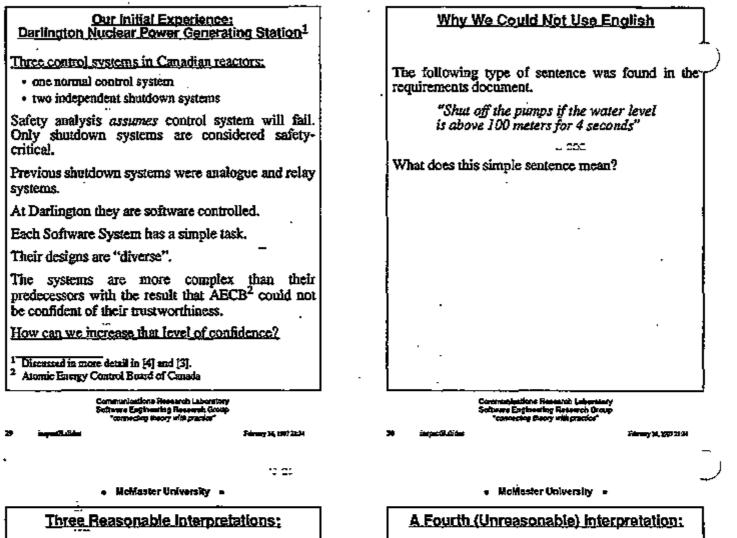
The definition of "well-structured" should not be based on the absence or presence of certain control structures. It has to do with the ease of decomposition. [2]

> Communications Research Laboratory Software Engineering Research Group formating theory with practice?

Person 28, 200 23 24

Ferrary 34, 100721-34





"Shut off the pumps if the mean water confevel over the past 4 seconds was above 100 meters".

320

ing on Labor

 $\left[\left(\int_{7-4}^{T} WL(t) dt \right) + 4 > 100 \right]_{300 \text{ SV}}$

"Shut off the pumps if the median water level over the past 4 seconds was above 100 meters".

 $(MAX_{(t-4,t)} (WL(t)) + MIN_{(t-4,t)} (WL(t))) + 2 > 100$

"Shut off the pumps if the "rms" water level over the past 4 seconds was above 100 meters".

Communications Security Laboratory Software Engineering Research Group "connecting theory with practice"

Tenney 36, 1997 22:34

Sottowere Engineering Research + *conversion theory with practic

"Shut off pumps if the minimum water level over the past 4 seconds

 $MiN_{(T-4,3)}[WL(t)] > 100$

It is a disaster waiting to happen!

If you use natural languages, there are thousands of

was above 100 meters".

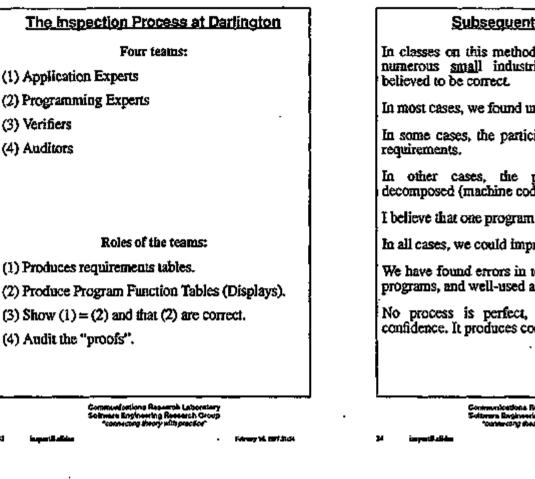
This is the most literal interpretation!

such phrases waiting to "bug" you.

McMaster University

32 impact@ulit

Manay 34, 19972104



Essential Point: Divide and Conquer

The initial decomposition is essential. Attempts to simply scrutinise the program fail.

Trying to read the program the way a computer would is much less effective. Logically connected parts may be far apart.

The use of tables is essential. It breaks things down into simple cases so that

- We can be sure that all cases are covered
- Each case is straightforward.

We consider all variables, but one at a time.

We consider all cases, one at a time.

We can take "breaks", go home and sleep, even take holidays, without losing our place.

Using displays and tabular summaries is far more work than Fagan's English paraphrasing, but it imposes a discipline that helps.

Communications Research Laboratory Solimane Engineering Research Group "connecting theory with practice"

Instant dirity

Fairway 16, 1997 20,04

Commentations Research Laboratory Software Engineering Research Group "commence theory with practice"

ionentii siide

W16.1072108

McMaster University

Subsequent Experience

In classes on this method, we have applied this to numerous small industrial programs that were

In most cases, we found unexpected errors.

In some cases, the participants could not state the

In other cases, the program could not be decomposed (machine code w/o documentation).

I believe that one program was correct.

In all cases, we could improve the program.

We have found errors in textbook programs, library programs, and well-used and tested programs.

No process is perfect, but this one engenders confidence. It produces code that people trust.

Communications Research Laboratory Software Beginsering Research Group "contenceng sheary with practice"

McMaster University

The Other Essential Point: Precise, Abstract Descriptions

Having lots of little parts is not enough.

We have to be sure that the parts fit together.

We have to be able to do that without page-flipping.

Each part's behaviour must be precisely summarised without giving intermediate states.

We must be sure that the description at the bottom of one display will be identical with that at the top of another display.

These global checks can, and have been, mechanised.

Precise descriptions are painstaking work, but if quality is important, they are essential,

It's not always easy!

The most critical step, besides decomposition, is finding a good representation for the state space.

A 1:1 relation between names and elements of the data structure cannot be assumed.

When preparing the displays, the creative step is data state representation.

earch Labora

wy I C 2007 21:34

en R

- **1**

. Transference

Some Suggested Reading

- Parnas, D. L., Weiss, D. M., "Active Desig Reviews: Principles and Practices", Proceedings of the 8th International Conference on Software Engineering, London, August 1985.
 Also in Journal of Systems and Software, December 1987.
- (2) Parnas, D. L., Madey, J., Iglewski, M.,
 "Precise Documentation of Well-Structured Programs",
 IEEE Transactions on Software Engineering, Vol. 20, No. 12, December 1994, pp. 948 - 976.

ļ

- (3) Parnas, D. L. "Inspection of Safety Critical Software using Function Tables", Proceedings of IFIP World Congress 1994, Volume III, August 1994, pp. 270 - 277.
- (4) Parnas, D. L., Asmis, G.J.K., Madey, J., "Assessment of Safety-Critical Software in Nuclear Power Plans", *Nuclear Safety*, vol. 32, no. 2, April-June 1991, pp. 189-198.

*			
Construction			
	Englis and by		
			- crange
*	inding photos	and a state	

Los y 14, 267 21 34

i

and the second

The Problem of the Dutch national flag¹

There is a data type color ≝ {blue,red,white} There is an abstract data type "buckets".

Variables of this type may be used as a vector of N "pebbles" of "color" type, where $N \ge 0$ is an integer.

The only operations on v are: PUT(i,c), LOOK(i), SWAP(i,j)

Design a procedure to rearrange (if necessary) the pebbles in the order of the Dutch national flag using no Arrays, and calling LOOK(i) once for each value of i.

¹ Introduced and (perhaps) solved by E. W. Dijkstra in 1976

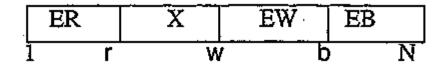
Communications Research Laboratory Software Engineering Research Group "connecting theory with practice"

24 nevercut.slides

7/3/96:1226

$1 \leq \mathbf{k} < \mathbf{r}$:	the k th bucket is in zone ER (number
	of buckets $r-1 \ge 0$)
$r \leq k \leq w$:	the k th bucket is in zone X (number
	of buckets wr+1 ≥ 0)
$W < k \le b$:	the k th bucket is in zone EW (number
,	of buckets $b - w \ge 0$)
$b < k \le N$:	the k th bucket is in zone EB (number
	of buckets $N-b \ge 0$)

This can be illustrated by the following figure.



Initially, r=1, and W=b=N, so that the zones ER, EW, and EB are empty. The program then proceeds by incrementing r, and decrementing W and b while making the necessary swaps, until the area marked "X" is empty because r = W+1.

Communications Research Laboratory Software Engineering Research Group "connecting theory with practice"

7/3/96:1226

```
program DutchNationalFlag (input, output);
 const
    N = 10;
 type
     color = (red, white, blue, blank);
    buckets = array [1..N] of color;
 var
     v : buckets;
     i : integer;
  function LOOK(i : integer) : color;
     begin
        LOOK := v[i]
     end;
 procedure PUT(i : integer; c : color);
     begin
        v(i) := c
     end;
 procedure SWAP(i, j : integer);
     var
        t : color;
     begin
        if ((i > N) \text{ or } (i < 1) \text{ or } (j > N) \text{ or } (j < 1)) then
              writeln ('wrong index passed to SWAP')
        else
           begin
               t := v[i];
               v[i] := v[j];
               v[j] := t
           end
     end; {SWAP}
```

Communications Research Laboratory Software Engineering Research Group "connecting theory with practice"

```
procedure Decrease(var r, w, b : integer);
    var
       colr, colw : color;
    begin
       colr := LOOK(r);
       while ((colr = red) and (r < w)) do
          begin
             r := r + 1;
             colr := LOOK(r)
          end:
       if (r < w) then
          begin
             {DecW}
             colw := LOOK(w);
             while ((colw = white) and ((r+1) < w)) do
                begin
                   w := w - 1;
                   colw := LOOK(w)
                end;
             case colw of
                red:
                        begin
                            SWAP(r, w); r := r + 1
                        end;
                white: w := w - 1;
                blue:
                        begin
                            SWAP(w, b); w := w - 1; b := b - 1;
                            SWAP(r, w)
                        end
             end
          end;
       case colr of
          red:
                 r := r + 1;
          white: w := w - 1;
          blue:
                  begin
                      SWAP (w, b); w := w - 1;
                    b := b - 1;
                  end
       end
    end; {Decrease}
```

Communications Research Laboratory Software Engineering Research Group "connecting theory with practice"

27

```
procedure Rearrange(var r, w, b : integer);
   begin
      while (w \ge r) do
         Decrease(r, w, b)
   end; {Rearrange}
procedure DutchFlag;
   var
      r, w, b < integer;
   begin
      r := 1;
      w := N;
      b := N;
      Rearrange(r, w, b)
   end; {DutchFlag}
                              (MAIN PROGRAM BODY)
begin
   {initialize the object v}
    DutchFlag;
```

end. (DutchNationalFlag)

Communications Research Laboratory. Software Engineering Research Group "connecting theory with practice"

28 nevercut.slides .

7/3/96:1226

LEXICON

A. Auxiliary functions

card: set \rightarrow integer card(s) $\stackrel{df}{=} 1$ s i (i.e. number of elements in the set s)

flag: buckets \rightarrow boolean

flag(v) # Br,b [partial_flag(v,r,r-1,b)]

partial flag: buckets \times integer \times integer \times integer \rightarrow boolean

 $\begin{array}{l} partial_flag(v,r,w,b) \stackrel{\text{def}}{=} (1 \leq r) \land (r-1 \leq w) \land (w \leq b) \land (b \leq N) \land \\ \\ \forall i \ (1 \leq i \leq N) \ [\ ((i < r) \Rightarrow (v_i = \text{red})) \land \\ \\ ((w < i \leq b) \Rightarrow (v_i = \text{white})) \land \\ \\ ((b < i) \Rightarrow (v_i = \text{blue})) \] \end{array}$

Note: v_i is defined in part C of this Lexicon.

same_colors: buckets × buckets → boolean

same_colors(v1,v2) 🖉

 $(card([i] (1 \le i \le N) \land (vI_i = red))) = card([i] (1 \le i \le N) \land (v2_i = red))) \land$ $(card([i] (1 \le i \le N) \land (vI_i = white))) = card([i] (1 \le i \le N) \land (v2_i = white))) \land$ $(card([i] (1 \le i \le N) \land (vI_i = blue))) = card([i] (1 \le i \le N) \land (v2_i = blue))))$

B. Pascal external definitions and declarations

const N = {ineral non-negative integer}
type color = (red, white, blue);
type buckets = {vector{N, color} - cf. part C of this Lexicon}
var v : buckets;
procedure LOOK(i : integer);
{cf. part C of this Lexicon}
procedure SWAP(i, j : integer);
{cf. part C of this Lexicon}

C. vector(n,elem) Module Interface Specification

(0) CHARACTERISTICS

- type specified: vector(n,elem)
- features: single-object, generic
- foreign types: elem, <integer>, <positive_integer>
- foreign types: n: <positive_integer>, elem

Communications Research Laboratory Software Engineering Research Group "connecting theory with practice"

29

McMaster University ...

(I) SYNTAX

ACCESS-PROGRAMS

Program Name	Arg#1	Arg#2	Value Type
LOOK	<integer>;V</integer>		eiem
PUT .	<integer>:V</integer>	elem:V	
ŚWAP	<integer>:V</integer>	<integer>:V</integer>	

(2) CANONICAL TRACES

canonical(T) \Leftrightarrow T = [PUT(i, e_i)]ⁿ _ = [PUT(i, _)]ⁿ_{i \equiv 1}

EQUIVALENT NOTATION FOR TRACES

1

Trace	Equivalent notation
vLOOK(i)~	٧Į

(3) EQUIVALENCES.

T.LOOK(î) ⇒ T

T.PUT(i, c) 🖘

Condition	Equivalence
(l≤i≤a)	~ %wrong_index%
l≤i≤n	T1_PUT(i,c).T2 where T=T1_PUT(i,x).T2

T.SWAP(i, j) =>

Condition		Equivalence
	≦j≤n))	Swrong_index%
	 (i < j)	T1.PUT(i,x).T2.PUT(j,y).T3 where T = T1.PUT(i,y).T2.PUT(j,x).T3
$(i \leq i \leq n) \land (i \leq j \leq n) \land i$	(i = j)	Т
	(i>j)	T1.PUT(j,x).T2.PUT(i,y).T3 where T = T1.PUT(j,y).T2.PUT(i,x).T3

(4) RETURN VALUES

Program Name	Argument No	Value
LOOK	Value	e where #0 = T1.PUT(#1,e).T2

Communications Research Laboratory Software Engineering Research Group "connecting theory with practice" ~

30

Precise Documentation of Well-Structured Programs

David Lorge Parnas, Jan Madey¹, Michal Iglewski²

Telecommunications Research Institute of Ontario (TRIO) CRL, McMaster University, Hamilton, Ontario, Canada L8S 4K1

ABSTRACT

This paper describes a new form of program documentation that is precise, systematic and readable. This documentation comprises a set of displays supplemented by a lexicon and an index. Each display presents a program fragment in such a way that its correctness can be examined without looking at any other display. Each display has three parts: (1) the specification of the program presented in the display, (2) the program itself, and (3) the specifications of programs invoked by this program. The displays are intended to be used by Software Engineers as a reference document during inspection and maintenance. This paper also introduces a specification technique that is a refinement of Mills' functional approach to program documentation and verification; programs are specified and described in tabular form.

1 Introduction

The process of program development has been thoroughly studied for nearly 30 years and useful insights have been gained. However, the focus of this work has been on designing the *first* version of a program. If a software product is successful, the program will have many more readers than writers and will be studied and revised many times. Moreover, while the writers have had the time to become closely familiar with the program, most readers will not have that huxury. We consider the needs of readers, e.g. reviewers and maintainers, to be at least as important as the needs of program designers. Although proper decomposition of the software into modules will reduce the complexity and length of programs, there will still be programs whose length makes them difficult to understand. This paper presents a method that can be used by developers to present their programs in a way that makes review and maintenance easier. The heart of the method is a way of precisely summarizing the effects of a program component, so that reviewers and maintainers do not have to study that code when looking at components that interact with it. The program and documentation are organized in such a way that the information needed to study a component is presented together with that component. This method is intended for programs that are well-structured in the sense defined later in this paper.

The present report is a revised version of [24]; it will appear in IEEE Transactions on Software Engineering,

1.1 On the role of documentation

Anyone who has ever seriously read a lengthy program produced by others (for example to inspect it or to make changes to it) realizes the importance of documentation. Some argue that well-written programs are self documentation. Practical experience suggests that this is true only for small programs; human beings cannot easily understand long programs. When asked to study such programs, we tend to focus on little details while making use of inaccurate

Permanent address: Institute of Informatics, Warsaw University, Banacha 2, 02-097 Warsaw, Poland

² Permanent address: Département d'informatique, Université du Québec à Hull, Hull, Québec, Canada J8X 3X7

descriptions of the overall structure. The combination of a large amount of detail with inaccurate or vague descriptions of the structure makes it quite common for serious errors to escape the reviewers' attention.

A design concept or algorithmic method that was obvious to the programmer at the three the program was written will not be obvious to other programmers, or even to the same programmer, one year later. Even if the program was developed using a systematic refinement process, there are few traces of that process in the final code. Although the program's author may have thought of the program in terms of a set of building blocks, each with a clearly defined function, it is not easy to identify those blocks and induce their functions by looking at the final code.

1.2 Studying long programs

When studying a long program, we must decompose it into small parts and then, provisionally, associate a function with each one. We must then convince ourselves of two things: (1) if each part implements its assigned function, the whole program will be correct, (2) that each part implements its assigned function. Frequently, we find that our provisional assumptions were not exactly what the programmer intended. Then, after revising our initial division and function descriptions, we try again. In principle, this iterative process converges and we learn whether or not the program is correct. In practice, we usually give up before we have a complete understanding of the program. The process terminates when we run out of time or patience.

1.3 Conventional documentation

Experienced development organizations have long recognized the need for documentation and there are extensive documentation standards. Unfortunately, when one tries to use this documentation, it is not found to be very useful. Often, the document includes a narrative description of the program - a translation of the program into a "natural" language. For people with an understanding of programming, it is usually easier to read the program itself than prose that attempts to say the same thing. Our natural languages were not intended to be used for precise descriptions where small details are critical. Most documentation encountered in industry is vague, inaccurate, and incomplete.

When documenting programs, there seems to be a tendency to focus on the details that we think will be hard to member while ignoring the basic structural decisions, which seem obvious. Later, readers find that the structure is obvious and the details are overwhelming. Moreover, most documentation is informally organized. Even when the desired information is present, it is not obvious where it will be found. When the information is found, it is often inconsistent or inaccurate. Industrial experience suggests that a huge portion of the "maintenance effort" goes into finding information and then finding an expert who can confirm or correct the information that was found.

The inadequacies of most software documentation can, in part, be blaned on the differences between standard engineering practice and the way that software systems are designed. In engineering, the production of design documents plays a key role - it is rare to find an engineer proceeding by building first and documenting later. In engineering, mathematics is extensively used to provide accurate and detailed descriptions of the products to be built; the need for precise descriptions of each component of larger products is almost universally accepted. In contrast, software systems are commonly produced *before* proper documentation is written; documentation is not viewed as a part of the design activity but as an additional task required by bureaucratic regulations or ignorant customers. The use of mathematics in describing programs is rare. As a result, the documentation is of limited value for programmers, reviewers and maintainers.

1.4 Design through documentation

The methods presented in this paper must be understood in the context of the complete documentation scheme described in [19].

It is widely accepted that the documentation of a computer system must include a software requirements docu-

ment (consisting of a system requirements document and a system design document). These documents provide a black-box description of the system as a whole, a description of the hardware structure, and a black-box description of the software. Detailed discussions of these documents can be found in [5, 6, 28, 29].

Because large software systems are seldom the product of a single person, the task of constructing them must be split into several smaller work assignments. Each assignment is to design and implement a group of one or more programs, which we call a module. In well-structured systems, the programs in a module share access to a private data structure and implement one or more abstract objects. We call programs that are part of the module, and can be used from outside the module, the access-programs of the module. Programs that are part of the modules never read directly from, or write directly to, the internal data structure of a module; they always use a module's access-programs to get information about, or change the state of, any objects created by that module [16]. We recommend a software module guide, which describes the structure of the software system by indicating the design decisions hidden in each one [23]. For each module identified in the module guide, there should be a module interface specification, which provides a black-box description of the behavior of the objects created by that module. Our approach to specification of module interfaces (the trace assertion method) is illustrated in Appendix B and described in [21, 3].

For every implementation of a module interface specification (there may be several), there should be a document describing the *module internal design*; that document must describe the internal data structures and the effect of the module's access-programs on the state of that structure. The contents of these documents are defined in [19], which contains a more general discussion of the role and structure of documentation in software engineering. Examples of a detailed software requirements document can be found in [5, 28].

This paper focuses on the documentation of programs within a module. The documentation described here complements the documents mentioned above.

1.5 The responsibilities of program designers and reviewers

We believe that the reviewer or maintainer of a program should never have to guess its structure. The iterative process described in Section 1.2 must be eliminated. Programs should be presented to the reviewer and maintainer as a collection of small parts, each with a precise description of its function. The structure should be explicitly and precisely described in the documentation. It should be possible to review the small parts separately and know that, if each of the components is correct, the whole program is correct. In other words, the decomposition phase of the review process should not be repeated by the reviewer; it should be communicated by the designer. The reviewers must check that the structure is a good one, but their primary responsibility should be checking each of the small fragments against the description of its function.

It is clear that we are asking more work from the designers than they usually do. We are asking them to write down, systematically, information that reviewers and maintainers would otherwise have to discover for themselves. Because there will be more readers than writers, and because the writer already knows the information, we believe that the combined cost of developing and maintaining the product will be lower if the writer presents the program as proposed in this paper and the documentation is kept live by revising it each time that the program is revised. Moreover, our experience suggests that the quality of the program will be improved as a result of requiring the programmer to produce the documentation.

1.6 The use of mathematics in documentation

Our method is based on a mathematical model of programs and uses mathematical notation to provide precise descriptions of programs. Although mathematics is not commonly used in programming practice, we believe that the ability to use mathematics in this way will be the hallmark of Software Engineers in the future.

Most demonstrations of the use of mathematical methods in software engineering emphasize program develop-

ment or verification. This paper focuses on documentation. While we believe in systematic development, we believe that the documentation delivered with a program should not depend on the program development process. This paper discusses the documentation that should be associated with a program, not the procedure for developing the program.

Many papers on formal methods for program development emphasize the idea of proving a program to be correct. Our paper is less ambitious. Although we believe that the mathematical documentation we describe could be used as input to a program verification process (our notation is close to classical predicate logic), our emphasis is on documentation that is valuable whether or not formal proof is attempted. We have used this type of documentation as input to an inspection process [22], but this paper does not discuss formal verification.

1.7 Introduction to the "Display Method"

This paper introduces a method of documenting well-structured programs called the Display Method. It requires designers and implementers to present their programs as sets of displays. The method is based on the well-known fact that a well-structured program can always be written as a short text in which the names of other programs³ may appear and the programs named can also be short. The down-side of such an organization is that there will be many programs and to understand any one of them, one must understand several others. We overcome this by presenting the material in displays. A *display* is a document in which a program is presented in such a way that its correctness can be examined without looking at other displays.

Though the Display Method can be used with any specification technique (and any imperative programming language), we decided to illustrate it using a refinement of Mills' approach⁴ [13, 14]. We have chosen to base our work on Mills' method, rather than approaches that are more popular, because we find it more suitable for large programs. Unlike Floyd [3], Hoare [7], Dijkstra [2], and their followers, Mills, although equally rigorous, does not include axiomatic descriptions of programming language statements among his basic definitions. Instead, he assumes that the cograms, from which other programs are constructed, can be described by mathematical functions. Since this asamption is valid for all deterministic programs, one can apply Mills' approach even when the component programs is quite long and complex. This allows the same method to be used for well-structured programs of any size.

Many other methods do not deal with the problem of how to assemble small programs into large ones. For examit, if one were to mimic the techniques used by With for the eight queens problem [27], one would keep repeating the parts of the text that were developed early in the refinement process. For a long program, this would not be practical. Program texts would grow so long that no one could keep them under full intellectual control. Other presentations of moderate-sized programs are confusing because it is not clear how the small sections fit together (cf. e.g. chapters 14 and 24 in [2]). Our method avoids both problems.

In documentation, the notation is very important; documents are to be read by experts from a variety of fields and Nould be easily understood. We must apply the principle of "divide and conquer" when designing notation; readers whould not have to parse iong expressions. Our approach is based on the use of tables to describe mathematical functions, relations, and sets [18]; such tabular notation has already been used in practice (e.g. in safety-critical software for a nuclear plant (22]) and has proven practical.

Some readers will observe that, in our examples, the volume of the documentation is much greater than the volume of program code. This is a consequence of the need to use small, but nontrivial, examples in a paper of this sort. The length and complexity of a precise description of a program's effect does not necessarily increase with the length of the program. In fact, it often happens that the description of the effects of a part of a program is more complex than the description of the whole program. Consequently, the ratio of program size to program documentation size is under

³ Note that these named programs need not be subroutines. In the text submitted to the compiler some of the program names may have been replaced by the text of the program itself.

⁴ Although Mills is the best known proponent of this approach, similar ideas were independently discovered by many others.

the control of the document's author. When documenting long, but easily understood, programs, it is not necessary to describe the behavior of small components; consequently, the ratio of code size to documentation size increases. In practice, the components identified will be longer than those in this paper's examples.

1.8 Organization of this paper

In the next section, we review some old issues about the structure of programs. Section 3 contains some basic definitions used in our approach to program description. Section 4 presents the main ideas of the Display Method and introduces important notational conventions. The method is illustrated by two complete examples (presented in appendices). A discussion of these examples and some sample displays are presented in Section 5. The lessons learned from previous experience with the proposed approach, and some future plans, are described in the final section.

2 Well-structured programs

This section motivates restricting the structure of programs, and then states the constraints proposed. While some researchers consider the themes in this section obvious, many practitioners continue to ignore them.

2.1 Hierarchical control structure in programs

The well known "structured programming" constructs, such as "while" and "if then else" have two very useful properties:

(1) programs constructed using them can be decomposed into a hierarchy of parts (with lower level parts completely contained in an upper level part) using simple parsers, those parsers need not even distinguish one identifier from another,

(2) the semantics of the total program can be determined from the semantics of its parts, using simple operations (cf. e.g. [17, 20]).

Further, the semantics of the program can be determined in a flexible sequence, finding the semantics of inner parts first and finding the semantics of a sequence of programs constructed using "," either left to right, right to left, or a mixture - as one prefers. In fact, the work need not be sequential. In contrast, the use of "go to" and labels makes it. difficult to find a decomposition in which the components have simple semantics.

The above properties are important because they make it easier to study a long program one small part at a time, and to do so without a previous understanding of the overall structure of that program. In contrast, when a program is constructed using labels and unrestricted jumps, considerable understanding of the program is needed in order to decompose it into parts that can be studied independently.

Programs having the desired properties are often referred to as having a *hierarchical control structure* or as wellstructured programs. The Display Method is intended to be used for such programs.

2.2 Use of data abstractions

Even the best structured program will be difficult to explain and understand if it is presented in terms of complex. data structures. Essential information about the nature of the data and algorithm can be obscured by representational details.

Complex data structures should be encapsulated (or hidden) by the introduction of new data types that have been designed specifically for the type of data being stored. Such specially designed data types, known as *abstract data types* (because they allow the reader to abstract from the actual representation of the data), were introduced into the literature by Dijkstra [1]. The principle of information hiding, long used by very good programmers, was first discussed explicitly in [15].

October, 1994

Precise documentation of a program that uses abstract data types is not possible unless the properties of the abstract operations are also precisely documented. In this paper we presume that the abstract types are implemented by modules whose properties have been specified by a module specification method such as that discussed in [21] or by one of the algebraic methods. However, the examples in this paper have been selected so that they can be understood without an understanding of module specifications.

2.3 Discipline vs. notation

It will be seen that the usability of the discipline proposed in this paper is independent of:

- (1) the notation used to present the information in a display,
- (2) the language used for coding the program, and
- (3) the method used to verify the displays.

The present paper focuses on the contents of the displays, using one programming language and one of many possible notations for presenting specifications. We have chosen the Pascal language [9] for the initial examples, not because it is ideal but because it is familiar. We have chosen to use tabular representations of LD-relations for reasons explained in Section 3, but we believe that the display method could be adapted for use with other notations such as VDM [10]. While we do not present formal verifications, we claim that the information necessary for verification of any display is contained in that display and the lexicon.

3 Mathematical description of program effects

In this section we show how to use standard mathematical concepts to describe the effect of program execution. We introduce the LD-relation [17, 20] and its application to program description and specification. Those wanting to use this method must read this section carefully. The literature contains many notations that are similar but differ from this one in subtle ways. In particular, the meaning of our notation is (necessarily) different from that of both Hehner [4] and VDM [10]⁵; confusion can arise if one assumes otherwise.

3.1 Finite state machine approach

A digital computer can usefully be viewed as a finite state machine. For our purposes such a machine is one that is always in one of a finite set of states and whose operation consists of a sequence of state-changes, i.e. transitions from state to state.

Definition 1:

We will use the term "program" to denote a description of state-change sequences in a finite state machine. Programs may describe both finite (terminating) and infinite (non-terminating) state-change sequences.

Let P be a program and let U be the set of states of a digital computer. The following terminology and notation will be used in the sequel:

Definition 2:

- A complete state-change sequence described by P is called an execution of P.
- The set of executions of P that begin with the state x, $(x \in U)$, is denoted by $e_P(x)$, and x is called the *starting state* of the sequences in that set. The set of all executions of P is denoted by Exec(P,U).

⁵ The work described in [4] stresses the description of programs by a single predicate, which limits the ability to provide complete descriptions of non-deterministic programs. VDM only describes the behavior of a program when started in states that satisfy a precondition that guarantees termination. We chose a method that allows complete description of any program.

- If there exists a finite execution in ep(x) with final element z, then:
 - we write $\langle x, ..., z \rangle \in e_p(x)$,
 - we say that this execution terminates (in z) and call z the final state (of this execution).
- If $\langle x, ..., z \rangle \in e_p(x)$, we also say that the program P may start in x and terminate in 2,
- If ep(x) contains an infinite sequence, we say that this is a non-terminating execution, and denote it by <x, ...>.
- If there exists a state x, (x ∈ U), such that e_p(x) contains two or more distinct executions, then P is a non-deterministic program.
- If for a given state x, (x ∈ U), every member of e_p(x) terminates, then x is called a safe state for P. The set of safe states for P is denoted by S_P.

۵

D

Ð

3.2 Limited-domain relations (LD-relations)

If we are not interested in the intermediate states of executions, then every deterministic program can be described by a *program function*, a function whose domain is the set of safe states and whose range is the set of final states [13]. Non-deterministic programs cannot be fully described by program functions. First, a program started in a safe state may terminate in one of several distinct final states; thus a relation must be used and not a function. Second, a program started in a state that is not a safe one may sometimes terminate and sometimes not; a relation on the set of states does not provide sufficient information to distinguish between safe and unsafe states.

In [17] one possible solution⁶ to the latter problem was suggested: instead of a relation we use a pair, (relation, set). This set will be used to provide the necessary additional information. The definitions that follow describe this solution. We begin by defining some formal structures, and describe how these can be used to describe and specify programs.

Definition 3:

- A binary relation R on a given set U is a set of ordered pairs with both elements from U, i.e. R ⊆ U × U. The set U is called the Universe.
- The set of pairs R could also be defined by its characteristic predicate, R(p,q), i.e. R = ((p,q):U×U | R(p,q)).

The domain and the range of R can be expressed as follows:

 $Dom(R) = \{p \mid \exists q (R(p,q))\}, \qquad Range(R) \Rightarrow [q \mid \exists p (R(p,q))]\}.$

In the sequel the term "relation" means "binary relation".

Definition 4:

Let U be a set. A limited-domain relation (LD-relation) on U is an ordered pair $L = (R_{Ls} C_L)$, where:

- R_L , the relational component of L, is a relation on U, i.e. $R_L \subseteq U \times U$,
- C_L , the competence set of L, is a subset of the domain of R_L , i.e. $C_L \subseteq Dom(R_L)$.

3.3 Applications of LD-relations

An LD-relation can be used both to specify and to describe programs. A program specification is a statement of the requirements that an acceptable program must satisfy. A program description is a representation of the visible behavior of a specific program. A specification may allow behavior that is not actually exhibited by the program. Since the same mathematical structure is used for both descriptions and specifications, each must be labelled to indicate the intended interpretation of the information. The following sections explain our usage of these terms more precisely.

⁶Other, mathematically equivalent, approaches introduce a special symbol to represent non-termination, cf. e.g. [12]. The approach chosen here allows representation in terms of variable values without the addition of any special symbols or states.

3.3.1 Program descriptions

As was mentioned in Section 3.2 a deterministic program can be described by a program function. We can generalize this notion, as follows:

Definition 5:

- Let P be a program, let U be a set of states, and let $L_P \neq (R_P, C_P)$ be an LD-relation on U such, that:
 - $(x, y) \in \mathbb{R}_{p} \Leftrightarrow \langle x, \dots, y \rangle \in \operatorname{Exec}(\mathbb{P}, \mathbb{U}),$

 L_p is called the *LD*-relation of *P* and the description of *P*.

If C_p = Dom(R_p), then (by convention) the competence set need not be given explicitly. In other words, if C_p is not given, then it is, by default, Dom(R_p).

비고

One should note the following consequences of this definition:

- if $x \in C_p$, P always terminates when started in x and if $(x, y) \in R_p$, P may terminate in y,
- if x ∈ (Dom(R_p) C_p), the termination of P when started in x is non-deterministic; in that case if (x, y) ∈ R_p,
 P may terminate in y, but it might not terminate at all,
- if $x \in \text{Dom}(R_p)$ and P starts in x, then P will never terminate.
- If P is a deterministic program, then the relational component, R_p, is a function, C_p = Dom(R_p), and hence L_p is the program function defined in [13]. Hence, our approach is "upward compatible" with that of Mills.

3.3.2 Specification of programs

We can also use LD-relations to specify a program. In the general case one may be given a set of LD-relations and be asked to write a program that satisfies at least one of them.

Definition 6:

Let $L_p = (R_p, C_p)$ be the LD-relation of a program P (where U is the set of states). Let S, called a *specification*, be a set of LD-relations on U, and let $L_s = (R_s) C \beta$ be an element of S. We say that:

² P satisfies the LD-relation L_s , iff $C_s \subseteq C_p$ and $R_p \subseteq R_s$,

- P satisfies the specification S, iff P satisfies at least one element of S.

Often, S has only one element. If S is a specification and $S = \{L_S\}$, then we can also call L_S a specification. This is the usual case and the only one illustrated in this paper.

If L_p is used as a specification, P will satisfy it. However, P will satisfy many other specifications and other programs may satisfy L_p.

4 The Display Method of program documentation

In the Display Method, program documentation consists of a set of *displays*, supplemented by a *lexicon* and an *index*. This section explains these concepts.

4.1 Displays

A well-structured program can usually be written as a short text in which names of other programs may appear.

⁷ Please note that C_p is not the same as the precondition used in VDM [10] and other methods. LD-relations provide a complete description of the behavior of a program, not just a description of its behavior when the starting state is in C_p. R_p is a description of the behavior within its domain, not just within C_p.

These named programs can also be short and can include the names of other programs. By a *display* we mean a concise document, (preferably 1-2 pages), in which a short program is presented in such a way, that its correctness can be determined without examining other displays. More precisely:

Definition 7:

A display is a document that consists of the following three parts:

- P1: a specification for the program presented in this display,
- P2: the program itself. The names of other programs may appear in this text; we say that the these programs are invoked in this display,
- P3: specifications of all programs (other than that specified in P1⁸) invoked in P2 that are not known⁹.

Note, that the terms "program" and "invocation" are to be understood in a generic sense. A name appearing in the program P2 may represent a procedure call (in which case it will usually be followed by actual parameters) but may also be treated as a macro call, to be replaced by a sequence of instructions. In either case, the construction of the resulting program by merging the P2 parts of all displays should be a simple operation that can be done automatically. As discussed below (cf. Section 4.4), if an invoked program is not an available¹⁰ program, its specification must appear as P1 in another display.

2:

4.2 The lexicon

To avoid repetition of information in several displays, and the maintenance problems that result from redundant information, we place that information in a separate document, called the lexicon.

Definition 8:

A *lexicon* is a dictionary containing definitions of terms used in the program being documented. It will contain the definitions of any mathematical functions, programs constants, types, etc. that are used in more than one display.

We refer readers to the lexicon wherever the information that it comains would have appeared.

4.3 The index-

To help those studying a program we also recommend an index.

Definition 9:

An index is a list of all the variables, programs, etc. indicating where those items appear in the displays. If some names are used with more than one meaning, we also describe the category of each name,

D

D

4.4 Completeness and correctness

Each display can be reviewed without any reference to other displays; its correctness can be verified without looking at the implementation of either the programs that are invoked in that display or the programs that invoke the program it describes.

⁸ Note that if a program invokes itself recursively, one should not include the specification of that program in its own P3.

⁹ A known program is one that does not require a specification. The semantics of known programs are assumed to be understood. Every project should have a list of programs that are considered to be known.

¹⁰ An available program is one that exists in a project or system library. We need not have a display for an available program. Available programs are not necessarily known programs. Known programs are usually, but not always, available.

Definition 10:

- A display is correct if the program in P2 will satisfy the specification in P1, provided that the programs invoked in P2 satisfy the specifications given in P3.
- A set of displays is complete if, for each specification of a program (except an available program) that is found in P3 of a display, there exists another display in which this specification is in P1¹¹.
- A set of displays is correct if (1) the set of displays is complete, and (2) all displays are correct.

A display can be supplemented by an additional part, P4, that contains a demonstration of its correctness. This could be either a description of the informal reasoning routinely done by a programmer, or a more formal argument. The existence of this additional section would make the reviewer's task simpler - one would not have to invent a "proof", only to check one. In the present paper we do not supply P4.

4.5 Notation

In the examples of displays in this paper we will use LD-relations for program specifications and the Pascal language for programs. The LD-relations will be represented in a *tabular* form [18]. The basis of such representation is the fact that every relation can be understood as a set of ordered pairs defined by its characteristic predicate (cf. Definition 3, Section 3.2). A predicate is also used to represent the competence set of an LD-relation.

4.5.1 Introductory conventions

This section introduces some useful notational conventions. It is usual to describe predicates using boolean expressions. The tabular notation used in the present paper will be explained by means of examples.

Convention 1:

Let P be a program specified by an LD-relation L = (R, C), and let $(v_1, ..., v_k)$ be the variables in P that constitute its data structure, v. Then:

- " "v₁" (to be read "v₁ before") denotes the value of the program variable v₁ before an execution of P,
- " vi" " (to be read "vi after") denotes the value of the program variable vi after a terminating execution of P,

• "'v " (to be read "v before") denotes the value of the data structure v before an execution of P.

- " v' " (to be read "v after") denotes the value of the data structure valuer a terminating execution of P.

Each pair in R will be of the form (' v_i , v_i '). Note that ' v_i and v_i ', as mathematical variables, could have been replaced in the definition of R by other symbols, but we would then have to establish an explicit correspondence between those symbols and the components of program data structure. Our notational convention makes the correspondence implicit in the variable names,

Convention 2:

If it is clear from the context that the programming variables are a, b, c, ..., then one may write "R(.)" instead of "R(('a, 'b, 'c, ...), (a', b', c', ...))".

Convention 3:

In examples we will often need to express the fact that some variables do not change their values during the execution of a program. We found it useful to introduce a predicate symbol NC ("Not Changed").

 $\mathsf{NC}(\mathsf{v}_1,...,\mathsf{v}_k) \Leftrightarrow (\mathsf{v}_1^{*} = {}^{*}\mathsf{v}_1) \land \ldots \land (\mathsf{v}_k^{*} = {}^{*}\mathsf{v}_k)$

Ω

Ð

¹¹ Note that completeness of the set of displays can easily be checked mechanically.

Convention 4:

When we write a boolean expression to characterize a set of program variable values, we always assume that programming variables can only have values appropriate to their types and do not state those restrictions explicitly.

Convention 5:

The variables that form the domain and range for a given LD-relation can be listed in the heading preceding the LD-relation and need not be repeated in the characteristic predicates.

4.5.2 Tabular representations

To explain the tabular notation used in this paper, we introduce the following simple problem:

PROBLEM

Write a program which finds the maximum of two integer values stored in programming variables.

Discussion:

The data structure of this program will consist of three variables of integer¹² type named a, b, and max. The first two will be used to store the input values, while the third one will store the result. We will require that the final values of a and b be the same as the initial ones. Note, that the initial value of max (i.e. 'max) is irrelevant.

The above considerations lead to the following specification of this program by an LD-relation, $L_s = (R_s, C_s)$:

$$- R_{s}(,) = \{ (a' = 'a) \land (b' = 'b) \land ((('a \le 'b) \land (max' = 'b)) \lor (('a \ge 'b) \land (max' = 'a))) \} \},$$

$$-C_{s} = Dom(R_{s}).$$

Tabular form:

The characteristic predicate of the relation Rs can be given in tabular form.

A direct representation of this predicate as a table, is as follows:

ฑ		ʻa≲ʻb	ʻa≥ʻb
a' '		'a	'a
р,	=	ъ	'Þ
max'	-	ð	'a

For ease of checking tables, we usually require that conditions that head columns be mutually exclusive¹³. In this case we should replace "\$" by "<", or "2" by ">". The first replacement leads to the following table:

T2	'a	≤ b	'a > 'o
a' :		'a	'a
b		5	Ъ
max" -		ъ	'a

¹² We will use different fonts to distinguish between programming language elements (e.g. "integer", "true"), and mathematical terms (e.g. "integer", "true").

• 4-

Ð

D

۰

¹³ This requirement is not strictly necessary, just useful. Eliminating heading overlap for tables that represent functions, cannot change their meaning and, consequently, does not result in overspecification. We show how to describe relations below.

The first two rows of TZ can easily be expressed conventionally. We can combine both notations as follows:

1	B		ʻa≤*b	'a > 'b	
	ភាax'	#	ъ	' B	

Using "NC" we can rewrite the above expression as follows:

B	Ç58,	ď < 8'	
max' 🛥	ъ	ʻa	∧NC(a,b)

• The conditions in 73 itself can be written in another way (which may make the table easier to read when expressions are long) - the string above a dotted line is treated as if it were repeated in each column below that line:

74	('a ≲		
	true ·	false	
max" =	Ъ	ťa	∧ NC(a, b)

- The conditions heading columns in 74 can be written in yet another form, as follows:

25	،		
	≤'ò_		
max" =	٦ ⁴	' a	^ NC(a, b)

 The equality operator in the "value after" phrase can be replaced by any other relational operator or by the vertical bar, "!". The latter is to be read "such that". When "i" is used, the entries in that row must be boolean expressions; the value of the variable must satisfy the predicate described in the relevant column. For instance, the row defining max' in the table II could have been written as follows: Note that the use of "i" allows the description of relations

Т	'a≤15	'a>′b	
max" [πax'≈ 'b	max' = 'a	~ NC(a, b)

or non-deterministic programs without having overlapping column headings.

The table identifiers: T1, T2, ... are optional and have no formal meaning.

0

4.6 Parameters and side-effects

Programs presented in displays will often use procedures. Procedures are not programs in the sense described above; they are program schemata, which cannot be described by functions or LD-relations. Procedures with formal parameters can be represented by program function schema, mappings from actual parameters to program functions, as described and illustrated in [8]. A procedure invocation, including the actual parameters, is a program in the sense of this paper. Here, we provide the program function corresponding to each actual invocation.

- (1) The specification of the procedure invocation will be written in terms of actual parameters. In the declaration of this procedure, however, formal parameters will be used. Both the specifications of invoked programs appearing in the declaration, and statements in the declaration body must be written in terms of the formal parameters of the procedure (and its other local or non-local objects). The binding of parameters is done according to semantics of the given programming language (Pascal, for the examples in Section 5).
- (2) For simplicity's sake, we will avoid any form of aliasing¹⁴ in our examples, e.g.:
 - · If more than one parameter is called by variable, then the actual parameters will be different variables.
 - If there are side-effects, then a variable external to the procedure body will not be passed as a parameter called by variable.

5 Examples

In this section we will illustrate the Display Method on two simple but complete examples written in Standard Pascal [9]. We decided to use simple and well-known problems to emphasize the main ideas of the proposed approach. The complete sets of displays with the lexicons and the indices are presented in appendices.

5.1 "Binary search"

We begin with a problem familiar to all programmers, so that we can focus on the display method.

5.1.1 Informal description of the problem

Given an integer x, and a list of $n \ge 1$ integers, a_1, \ldots, a_n in non-decreasing order:

- check whether x is among a_1, \ldots, a_n and return this information,
 - if x is among a_1, \ldots, a_n , find an index j such that $x = a_j$.

If the list is empty or not sorted, we require program termination but do not care what the program does because we assume that the program will not be invoked under such conditions¹⁵.

5.1.2 Discussion

- (1) A solution to this problem (by the well-known "binary search" method) will be presented as a Pascal procedure declaration and its invocation. It is the invocation that must satisfy the specification. This procedure declaration should be preceded by definitions and declarations of needed constants, types and variables, to set up the data structure whose values will form the state space.
- (2) The following assumptions are made about the correspondence between the description of the problem and Pascal programming language objects:
 - Integer numbers are represented by values of the standard type integer¹⁶.
 - The length of the list is represented by the constant n.
 - The list itself is represented by the value of the variable A of a type vector, defined as array[1...n] of integer.
 - The integer x is represented by the value of the variable x of type integer.
 - . The results are represented by the values of two variables: j of type integer, and present of type Boolean.

¹⁵ This is undoubtedly a foolish assumption in practice, but it is useful for illustrating the meaning of the notation. In this example, if a program is called when the assumptions are not satisfied, even the values of the variables x and A are allowed to change.

¹⁶ Recall that by convention we use different fonts to distinguish Pascal objects from mathematical ones.

¹⁴ Aliasing does not invalidate the basic theory or model used in our work. However, it complicates the representation of data states. In our examples, there is a 1:1 correspondence between identifiers and elements of the data structure at any point in the program. This allows us to represent state by a list of values in which each element corresponds to one identifier. If aliasing is allowed, or with dynamic data structures, one needs a more elaborate scheme for identifying data states.

- (3) We will specify, that:
 - The values of A and x should not change if the program is invoked under normal conditions.
 - If the desired index exists, then j will return its value and present will be true. If the index does not exist, present will be false and j can have any integer value.
- (4) The following observations and conventions are related to the data state:
 - Initially, the data state is determined by the values of the constant n and the variables A, x, j, and present.
 - The relational component R of the LD-relation should specify acceptable changes of these values (however constants, by definition, do not change and their values need not be mentioned).
 - For variables we will use the conventions introduced in the previous section.

5.1.3 Example of a display

We will present one display (the complete set is to be found in Appendix A). To help in understanding specifications, we begin by discussing P1 of this display in detail. We have numbered each line of part P1 and explain those lines in the notes below.

(1)	Find(x, A, j,	Find(x, A, j, present)				
(2)	R ₀ (,) = ((1 ≤	$R_0(\mathbf{s}) = ((1 \le n) \land \forall i [(1 \le i \le n) \Rightarrow (A_i) \le A_i]$				
(3)	$\exists i [(1 \le i \le n) \land (A[i] = X)] =$					
(4)		іпне	faise			
(5)	Ϊ I	*A[j'] = 'x	συς			
രി	present' =	true	false	∧ NC(x, A)		
		· · · · · · · · · · · · · · · · · · ·		-		

Specification

Notes on P1:

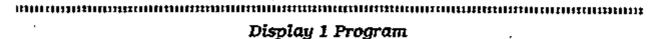
- (1) The procedure invocation "FIND(X,A,j,present)" lists actual parameters which form the data structure. If external¹⁷ variables were used, they need to be listed in this line.
- (2) Since the elements of the data structure are listed in line 1, we do not need to repeat them (Convention 2, Section 4). Without that convention we would have to write "R₀(('x, 'A, 'j, 'present), (x', A', j', present))" instead of "R₀(,)". Next note, that the expression "((1 ≤ n) ∧ ∀i [(1 ≤ i < n) ⇒ ('A[i] ≤ 'A[i+1])])" is true if the input sequence is non-decreasingly ordered.
- (3,4)This and the next line could have been written as one entry but we would have to repeat the long expression twice.
- (5) The phrase "J 1 mue" expresses the fact that the program will satisfy the specification no matter what the value of j is when the program terminates.
- (6) Notice that the logical values written here are Pascal constants. The other "*rue*" and "*false*" were mathematical constants. The phrase "NC(x, A)" expresses the requirement that the variables with input values remain unchanged.

In P3 of the display, the rows for low and high are not strictly necessary because the new values of those variables are not constrained. Since these tables represent the characteristic predicate of the relation, variables that are not mentioned are not constrained. We sometimes include such rows to make this more explicit.

¹⁷ We will use the term *external* to denote objects that are not local to a given program.

Display 1 Specification

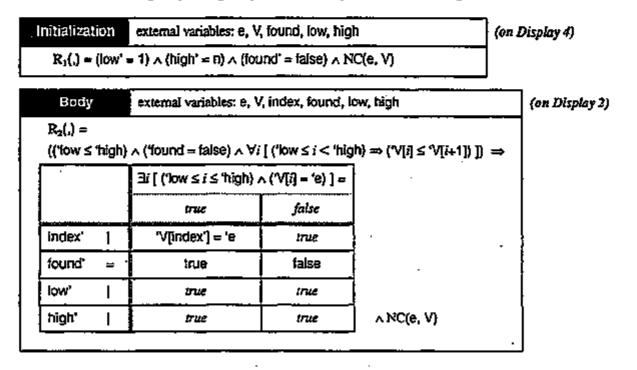
$R_0(.) = ((1 \leq$	$n) \wedge \forall i [(1 \le i < n) :$	⇒ ('A[i] ≤ 'A[i	+1]}]} ⇔
	∃i[(1≤i≤n)∧('	A[i] = 'x}] =]
-	irue	false	ĺ
ĵ	'A[]'] = 'x	true	1
present" =	true	false	^ NC(X, A)



Procedure declaration:

procedure Find(e : integer; V : vector; var index : integer; var found : Boolean); var low, high : integer; begin Initialization; Body end {Find}

Display 1 Specifications of Invoked Programs



October, 1994

CRL Report No. 295

END OF DISPLAY 1

5.2 "Dutch national flag" example

This example is based on [2], chapter 14.

5.2.1 Informal description of the problem

- (1) There is an abstract data type "bockets". A value of this type may be used as a vector of N elements of type "color", where N ≥ 0 is a fixed integer, and color floor (blue, red, white). Each element is called a "pebble" by Dijkstra. We introduce a variable of type buckets, v, c of type color, and i, j of type integer. The operations on v are:
 - PUT(i,c), which sets the value of ith element of v to c, if N>0, (i.e. puts the c-colored pebble into the ith bucket) and does nothing if N=0 or i is out of range.
 - LOOK(i), which returns the color of the pebble in the ith bucket and does nothing if i is out of range.
 - SWAP(i,j), which swaps pebbles between the ith and jth bucket, if i=j, and does nothing if i and j are equal or the arguments are out of range.
- (2) The type buckets and the operations PUT, LOOK and SWAP are defined more formally in Appendix B (in the lexicon) by a parameterized module interface specification using the trace assertion method [21, 8]. The initial value of v is assumed to be set externally.
- (3) We want to design a Pascal procedure that, given any initial arrangements of pebbles in v, "will rearrange (if necessary) the pebbles in the order of the Dutch national flag, i.e. in order from low to high bucket number first the red, then the white, and finally the blue pebbles." [2]. This procedure should:
 - cope with all possible special cases, including missing colors and N=0,
 - not introduce arrays of any sort, only a fixed number of variables of type integer and color, and
 - not use the operation LOOK(i) more than once for each value of i.

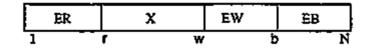
5.2.2 Discussion

Our solution (and the description in this section) is based on the original proposal by Dijkstra. We will assume the existence of the external Pascal variable v of type buckets, as presented in the problem description above, and that the Pascal procedures PUT, LOOK, and SWAP are both available and known.

Although the pebbles are of only three different colors, the fact that we can only inspect pebbles one at a time, together with the requirement that we can only inspect each pebble once, implies that throughout the arrangement process, we have to distinguish between pebbles of four different categories, viz. *established red* (ER), *established white* (EW), *established blue* (EB), and *as yet uninspected* (X). We will divide the row of buckets into four (possibly empty) zones of consecutively numbered buckets, each zone being reserved for pebbles of a specific category. For keeping track of the place of the zone boundaries we will use three integer variables, r, w, b, with the meanings:

> $1 \le k < r$: the kth backet is in zone ER (number of backets $r-1 \ge 0$) $r \le k \le w$: the kth backet is in zone X (number of backets $w-r+1 \ge 0$) $w < k \le 0$: the kth backet is in zone EW (number of backets $b-w \ge 0$) $b < k \le N$: the kth backet is in zone EB (number of backets $N-b \ge 0$)

This is illustrated by the following figure:



CRL Report No. 295

Initially, r=1, and w = b = N, so that the zones ER, EW, and EB are empty. The program then proceeds by incrementing r, and decrementing w and b while making the necessary swaps, until the area marked "X" is empty because r = w+1.

5.2.3 Example of a display

The complete set of displays including the lexicon and index is to be found in Appendix B. In the display below there are three auxiliary functions (predicates) used: flag, partial_flag, and same_colors. Their formal definition is given in the lexicon. Intuitively, flag(v) is true if the colors in v form the required final configuration (zone X is empty); partial_flag(v,r,w,b) is true if colors are grouped as on the above figure. The predicate same_colors(x,y) is true if x and y have the same number of red, white, and blue pebbles.

DISPLAY 1

Display 1 Specification

DutchFlag external variable: v $R_0(,) = flag(v') \land same_colors('v,v')$

Procedure declaration:

*** procedure DutchFlag; var r, w, b : Integer; begin r := 1; w := N; b := N; Rearrange(r, w, b) end {DutchFlag}

Display 1 Specifications of Invoked Programs

$$\begin{array}{l} \text{Rearrange}(\mathbf{r}, \mathbf{W}, \mathbf{D}) & \text{external variable: } \mathbf{V} \\ \mathbf{R}_1(.) \simeq ((\mathsf{'T} = 1) \land (\mathsf{'W} = \mathsf{N}) \land (\mathsf{'b} = \mathsf{N})) \\ & \Rightarrow \\ & (partial_flag(\mathsf{V}', \mathbf{f}', \mathsf{W}', \mathsf{D}') \land (\mathsf{W}' = \mathsf{f}' - 1) \land same_colors(\mathsf{'V}, \mathsf{V}')) \end{array}$$

END OF DISPLAY 1

6 Experience

The ideas reported in this paper are motivated more by practical experience than by theory. The theory has been introduced only to the extent that it was needed to provide a precise meaning for the notation. We have all had the frustrating experience of trying to read the mind of a programmer when trying to correct a program. The proposals in this paper represent our thoughts about what the programmer should have given us.

The method described in this paper is an improved version of the technique used in the inspection of safety-critical software for the Darlington (Ontario) Nuclear Power Generation Station [22]. It is important to understand that the Darlington experience was not an experiment; we did not gather data or make scientific observations. There was a job to be done and it had to be done as quickly as safety considerations would permit.

At the Darlington station, two safety-critical systems were, for the first time, implemented in software. The Atomic Energy Control Board of Canada (AECB) was not willing to allow the plant to operate until they were convinced of the correctness of the programs. Delays were very expensive for the owners of the plant, Ontario Hydro. The software had been ready for several years (because the rest of the plant was even further behind schedule), had been tested thoroughly, and was considered by its owners to be safe to use. However, the usual informal approaches to inspection did not provide the confidence level demanded by the AECB. The code, while not huge¹⁸, was sufficiently complex that the engineers who inspected it using informal methods could not be confident that they had considered all of the possibilities and found all of the errors.

One of the preliminary inspections demonstrated that the requirements documentation was not complete or precise. An error caused by misinterpretation of a sentence was discovered. As a result, the manufacturer was asked to produce a mathematical requirements document using [5] as a model. This document, which also used tabular representations of mathematical functions, was reviewed by nuclear safety experts.

It was also agreed that precise program documentation would be produced and used as the basis for an inspection procedure. Because the correctness of this code was considered vital to the safety of the plant, AECB, Ontario Hydro, and Atomic Energy of Canada Ltd. (AECL), were able to train approximately 60 engineers to produce and review tabular documentation. The inspectors had to identify program components and document them. The resulting tables were then used as the input to an open inspection process. Each table was presented to a review group and the anthors had to demonstrate that it was a correct description of the code. Generally, this involved going through the table on a column-by column, row-by-row basis. The tabular organization was extremely valuable because it made it easy to take breaks (the process went on for months) without losing context or continuity.

In addition to demonstrating that the tabular documentation of the programs accurately described the code, it was necessary to demonstrate that the tables describing the code described behavior that satisfied the requirements represented by tables in the requirements document. Generally, this involved a step-by-step transformation of one table until it matched the corresponding table in the other requirements document. The transformations were not mechanical; their correctness depended on properties of the functions used in the expressions and required human insight. Again, the tabular organization proved essential to allowing human beings with finite attention spans to compare two very detailed documents

In the Darlington work the documentation was not formally organized into displays. This led to a lot of page flip-

October, 1994

¹⁸ While line-counts are notoriously subjective, an outside expert ([11]) estimates the programs as containing about 2500 lines of FORTRAN and Pascal, plus about the same amount of code in assembler.

ping during the inspection process. Technological limitations also prevented us from using some of the notation in this paper. The work was done without the precise definitions in this paper and demonstrated the need for those definitions. In the Darlington work, for example, we did not use quantifiers and this led to problems when dealing with arrays in the program.

The methods described in this paper result from our reflection on the Darlington experience. The notations used here are the ones that we now believe we should have used in Darlington. The notation presented here has been used in more academic experiments including work done at Warsaw University and at McMaster University. Our conclusions are supported by experience gained when the Display Method was applied to examples larger than those presented in this paper (e.g. a simple data base) and implemented in different programming languages (Sun Pascal, Turbo Pascal, FORTRAN, C), cf. [26]. One interesting aspect of this McMaster University work was that it was done by an undergraduate with no prior exposure to formal methods or mathematical logic. He was able to document and repair a FORTRAN program that had been frustrating its owners in their attempts to repair it for many months. Our success did not surprise us, but it surprised the owners of the FORTRAN program who had reluctantly concluded that the program could not be salvaged.

The extensive experience gained in the Darlington work, and in subsequent uses of the method, has revealed where users of these ideas spend their time. We have found that much of the Engineer's time was spent on tasks that could be done by relatively simple tools. This work has led to tool projects at McMaster University, the Université du Québec à Hull, and Warsaw University, which will be described in the next Section.

7 Concluding remarks

We base this method on a very simple idea. Programs can only be understood in small chanks, so they should always be presented in small pieces. Each presentation must be complete in itself so that it can be studied without looking at the others. However, one can not follow this simple precept without finding a way to express the connections between the small sections. It does no good to have a collection of small programs, each one of them correct, if they do not fit together to make a large correct program. This observation led us to use a relational/functional model, both to specify the requirements that a program must meet, and to describe the behavior of a given program. While we found that conventional mathematical concepts were theoretically sufficient to describe these relations, conventional notation resulted in complex expressions that were hard to parse and understand. This led us to introduce a tabular notation that allowed us to describe the programs in a more readable manner. Without this notational progress, the original simple idea would not be as practical.

We began our work on the assumption that we were studying a method of program presentation. It soon became clear that the method was also a way of developing programs. Programs that had been developed before we began to document them, were found to have defects that became obvious when we started to present them in displays. Documenting programs using the display method can result in significant improvements in the quality of the program.

One advantage of this method is that one can speed up a review by employing more reviewers. The displays do not have to be reviewed in any special order and can be reviewed in parallel because they are independent. Even more important, if an error is found in Part 2 of a display, that part can be changed without necessitating modifications to any other displays unless Part 3 is changed. If we do find it necessary to change Part 3 of a display, other displays will have to be changed but we will know exactly which displays must be revised and checked.

The package of ideas that we have presented has proven valuable, but we believe that tool support is needed to make it practical for "everyday" programs. With current tools, it takes an excessive amount of effort to make sure that our expressions are syntactically correct and to achieve neat formatting. Moreover, it requires a high degree of discipline to perform simple checks on the displays, and to make sure that the specifications that are "copied" from the bottom of one display to the top of another are, and remain, identical. Checking lexicon entries requires annoying page-flipping or frustrating delays on the screen. Assembling the program segments to produce executable code by

hand is also a time-consuming process in which it is easy to introduce careless errors.

We believe that the situation can be ameliorated by building a set of tools that are designed to support this method of program development and documentation. We envision a system in which the central window presents a display, and other windows provide the relevant lexicon entries. In such a system, the "copying" of the specifications would be automatic and it would be impossible to change one without changing the other. The system would be capable of performing a completeness check and would remind us of specifications that could be found in Part 3 of one display but were not yet developed as Part 1 of another. Checking correctness remains a task for humans. We now have a prototype of such a tool. Other tools would provide syntactic and semantic checks and help us to format the displays. Work on direct support of the Display Method is being carried out at both McMaster University and Warsaw University. At the Université du Québec à Hull editors to support other types of formal documentation have been completed.

A system of this sort would be extremely valuable for people who develop software and even more valuable for those who maintain software products. It would be valuable even without any verification capability, but a simple theorem prover would allow us to make basic checks on the tables. In the future, documentation in this style could be used as input to more sophisticated provers. The information necessary for verification is present in these documents.

Because the documentation is mathematical in nature, it can be used to support testing. The tabular representations can be converted to "oracles", i.e. programs that evaluate the results of tests. If a program is tested against programs generated from it's documentation, developers are more likely to keep the program and documentation consistent. Work of this sort is described in [25].

Tools to make it easier to produce tabular representations of functions and relations in any kind of documentation are being studied and developed at McMaster University.

If readers take the time to compare our presentation of the problem of the Dutch National Flag with Dijkstra's original proposal [2] they will see the benefit of our approach. Dijkstra's presentation, though very illuminating, is dangeronsly unclear. Although he shows great discipline in developing the small program fragments that are presented in the text, he relies on informal discussions in describe how these are to be assembled into a complete working program. Four essential lines of program text in our solution cannot be found in the program fragments in the original version. Three of these lines are implied by an easily overlooked English sentence in Dijkstra's discussion of the program development. The fourth covers a simple case that seems to have been overlooked because the complete program structure was never presented. We know of several occasions where readers have been asked to examine the original description of the algorithm and then assemble working Pascal programs. Some readers simply assembled Dijkstra's program fragments - producing programs that were not correct. Others noted the conflictons in the English text and produced correct programs. We consider Dijkstra's description to be unclear; some have argued that it is wrong¹⁹. While no method guarantees error-free programs, we believe that the use of the Display Method with careful reviews of each display, makes such errors much less likely.

The problem of the Dutch National Flag reveals one of the limitations of our specification method. LD-relations, like predicate transformers and pre/post conditions, are unable to express the fact that the program is only permitted to inspect the contents of a bucket once. Relational methods limit the final state of the program, but the number of "LOOK" operations that have been carried out is not reflected in the final state with the data structure given. The definition of the buckets abstraction could easily be modified to distinguish between inspected and minspected buckets, but this would be modifying the data structure ordy to make the specification easier.

The binary search example illustrates the subtle ways in which programming language restrictions can affect the documentation. In Display 2 we had to introduce "med" but, because we were using Pascal, this variable's declaration should have been included in Display 1. If we had been using Pascal's predecessor, Algol 60, the declaration could

5 2 1

¹⁹ Dijkstra advised against bothering to assemble the final program, apparently because there was no need to look at it.

have been made where it was needed and kept local to the block in which it was used.

Acknowledgements

Dr. Gordon Stuart, then a Ph.D. student at the University of Victoria, contributed to the early development of these ideas. Wm. Wadge belped to clarify the concepts of LD-relations. Graydon Saunders worked on a diagrammatic form of displays for a mythical programming language that we call DAD. The many employees and consultants of AECB, AECL, and Ontario Hydro who applied some of these ideas to the review of the software for the Darlington Nuclear Plant [22] helped us by showing what was needed to turn some academic conceptions into a more mature technology. Mr. P. Filip Sawicki's work on his MSc Thesis [26] was an important step in exploration of the Display Method. We are grateful to David Weiss and to the many people who were kind enough to offer us comments on earlier versions of this paper. An exemptiatingly detailed review by referee #4 helped us to clarify our explanations in many ways. The editor's final suggestions were unusually helpful.

This research was primarily funded by the Telecommunications Research Institute of Ontario (TRIO), Other support was provided by the Canadian Institute of Telecommunications Research (CITR), by the State Committee for Scientific Research in Poland (KBN), by Digital Equipment's European External Research Programme (EERP) and by the Natural Sciences and Engineering Research Council of Canada (NSERC).

References

- Dijkstra, E.W., "The Structure of the 'THE' Multiprogramming System", Communications of the ACM, Vol. 11, No. 5, May 1968, pp. 341-346.
- Dijkstra, E.W., Discipline of Programming, Prentice-Hall, 1976.
- Floyd, R.W., "Assigning Meanings to Programs", Proceedings of the Symposium of Applied Mathematics, Vol. 19, 1968. Also in: Schwartz, J.T. (editor), Mathematical Aspects of Computer Science, American Mathematical Society, 1967, pp. 19-32.
- Hehner, E.C.R., "Predicative Programming, Part 1", Communications of the ACM, Vol. 27, No. 2, February 1984, pp. 134-143.
- Heninger, K.L., Kallander, J., Parnas, D.L., Shore, J.E., "Software Requirements for the A-7E Aircraft", NRL Memorandum Report 3876, United States Naval Research Lab., Washington D.C., November 1978, 523 pp.
- Heninger, K.L., "Specifying Software Requirements for Complex Systems: New Techniques and their Application", IEEE Transactions Software Engineering, Vol. SE-6, No. 1, January 1980, pp. 2-13
- Hoare, C.A.R., "An Axiomatic Basis for Computer Programming", Communications of the ACM, Vol. 12, No. 10, October 1969, pp. 576-580.]
- Iglewski, M., Madey, J., Pamas, D.L., Kelly P. C., "Documentation Paradigms", CRL Report 270, McMaster University, CRL, Telecommunications Research Institute of Ontario (TRIO), Hamilton, Ontario, Canada; July 1993, 45 pp.
- Jensen, K., Wirth, N., "Pascal User Manual and Report", Lecture Notes in Computer Science, Vol. 18, New York, Springer-Verlag, 1974 (second corrected edition 1976).
- 10. Jones, C. B., Systematic Software Development Using VDM, Prentice-Hall, 1986.
- Leveson, N., Personal Communication, 10 September 1994.
- Majster-Cederbaom, M.E., "A Simple Relation Between Relational and Predicate Transformer Semantics for Nondeterministic Programs", Information Processing Letters, Vol. 11, No. 4,5, December 1980, pp. 190-192.
- Mills, H.D., "The New Math of Computer Programming", Communications of the ACM, Vol. 18, No. 1, January 1975, pp. 43-48.
- 14. Mills, H.D., "Function Semantics for Sequential Programs", Proceedings of the IFIP Congress 1980, North Hol-

October, 1994

CRL Report No. 295

land 1980, pp. 241-250.

- Pamas, D.L., "Information Distributions Aspects of Design Methodology", Proceedings of the IFIP Congress '71, Booklet TA-3, 1971, pp. 26-30.
- Pamas, D.L., "On the Criteria to be Used in Decomposing Systems into Modules", Communications of the ACM, Vol. 15, No. 12, December 1972, pp. 1053-1058.
- Parnas, D.L., "A Generalized Control Structure and Its Formal Definition", Communications of the ACM, Vol. 26, No. 8, August 1983, pp. 527-581.
- Parnas, D.L., "Tabular Representation of Relations", CRL Report 260, McMaster University, CRL, Telecommunications Research Institute of Ontario (TRIO), Hamilton, Ontario, Canada; October 1992, 17 pp. +
- Parnas, D.L., Madey, J., "Functional Documentation for Computer Systems Engineering. (Version 2)", CRL Report 237, McMaster University, CRL, Telecommunications Research Institute of Ontario (TRIO), Hamilton, Ontario, Canada; September 1991, 14 pp.
- Paraas, D.L., Wadge, W.W., "Less Restrictive Constructs for Structured Programs", Technical Report 86-186, Queen's, C&IS, Kingston, Ontario, Canada, October 1986, 16 pp.
- Parnas, D.L., Wang, Y., "The Trace Assertion Method of Module Interface Specification", Technical Report 89-261, Queen's, C&IS, Telecommunications Research Institute of Ontario (TRIO), Kingston, Ontario, Canada, October 1989, 39 pp. (Available from McMaster University).
- Parnas, D.L., Asmis, G.J.K., Madey, J., "Assessment of Safety-Critical Software in Nuclear Power Plants", Nuclear Safety, Vol. 32, No. 2, 1991, pp. 189-198.
- Parnas, D. L'Clemenis, P. C., Weiss, D. M., "The Modular Structure of Complex Systems", IEEE Transactions on Software Engineering, March 1985, Vol. SE-11 No. 3, pp. 259-266.
- Parnas, D.L., Madey, J., Iglewski, M., "Formal Documentation of Well-Structured Programs", CRL Report 259, McMaster University, CRL, Telecommunications Research Institute of Ontario (TRIO), Hamilton, Ontario, Canada; September 1992, 37 pp.
- Peters, D., Parnas, D. L. "Generating a Test Oracle from Program Documentation", published in Proceedings of the International Symposium on Software Testing and Analysis, August 17 - 19, 1994.
- Sawicki, P.F., "Analiza metody SPS(R) weryfikacji programow" (in Polish), ["An Analysis of the SPS(R) Method of Program Verification"], MSc Thesis, Warsaw University, Institute of Informatics, 1992, 124 pp.
- Wirth, N., "Program Development by Stepwise Refinement", Communications of the ACM, Vol. 14, No. 4, April 1971, pp. 221-227.
- van Schouwen, A. J., "The A-7 Requirements Model: Re-examination for Real-Time Systems and an Application to Monitoring Systems", Technical Report 90-276, Queen's, C&IS, Telecommunications Research Institute of Ontario (TRIO), Kingston, Ontario, Canada, May 1990, 93 pp. (Available from McMaster University).
- 2' Lan Schouwen, A. J., Parnas, D. L., Madey, J., "Documentation of Requirements for Computer Systems", presented at RE '93 IEEE International Symposium on Requirements Engineering, San Diego, CA, 4 - 6 January, 1993.

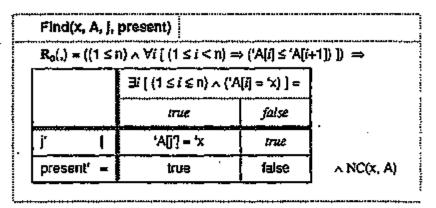
CRL Report No. 295

Appendix A

"Binary search" example presented on displays

The description of the problem and the discussion of the solution were given in Section 5.1. What follows is the formal documentation for the complete solution.

Display 1 Specification



Procedure declaration:

procedure Find(e : integer; V : vector; var index : integer; var found : Boolean); var low, high : integer; begin Initialization; Body end {Find}

Display 1 Specifications of Invoked Programs

nitializat	ion	external variables: e, V,	found, low, his	gh	(an Display 4)	
$R_1(.) = (iow' = 1) \land (high' = n) \land (found' = ialse) \land NC(e, V)$						
Body external variables: e, V, index, found, low, high					(on Display 2)	
$R_2(,) \Rightarrow$	"high)	\wedge ("found = false) $\wedge \forall i$:[("low≤i<"h	igh) ⇒ ('V[i] ≤ 'V[i+1]) :		
$\exists i [(`tow \le i \le `thigh) \land (`V[i] = `e)] =$					**	
		true.	false	-	1	
				-	1	
index'	l	*V[index*] = *e	irue			
index* found:	 =	"V[index"] = 'e true	true faise	-		
	 = }					

END OF DISPLAY 1

Display 2 Specification

Body	ļ	external variables: e, V	, index, found, k	ow, high	(from Display 1
R _z (,) = (("found	= falsi	e) ∧ ('low ≤ 'high) ∧ ∀i	[(fow $\leq i < highted for its initial constraints in the second $	$h) \Rightarrow (`V[i] \leq `V[i+1])]) =$	*
		$\exists i \; \{ \; (`kow \leq i \leq fuigb) \land (`V[i] = `e \} \; \} =$			
		true	false		-
index'	{	'V[រែវdex'] = 'e	truë		2 5 5 5 1 1
found"	=	true	faise		
low*		true.	true		
high'	1	true	true	∧ NC(e, V)	

Display 2 Program

New variable (to be declared in the embedding block): var med : integer;

Program statements:

(Body)

while not found and (low \leq high) do begin

med := (low + high) div 2;

Test

end

Display 2 Specifications of Invoked Programs

Test		external variab	(on Display			
$\mathbf{R}_{3}(.)=\langle$	'low ⊴	⊊'med ≤ fhigh)	**]
			V['med]]	
	Î	<'e	= 'e	>~e	- 	
index'	1	true	index' = 'med	true	~ -	
found	-	found	true	"tournd"	1	
low*	-	'med + 1	Yow	"low	1	1
high'	-	'high	'hîgb	'med - 1	A NC(e, V, med)	

END OF DISPLAY 2

October, 1994

Test		external variab	(from Display 2			
$\mathbb{R}_{\mathfrak{g}}(\mathbf{x}) = ($	low :	≲ 'med ≤ 'high)	⇒	+ +++++ = ++++ = = = = = = = = = = = =		
			"V["med]			
		<'a	= 'e	>'e	••• •	
index'	Ι	ţrue.	index" = 'med	17216		
found'	=	found	ນນອ	found		· ·
łow'	=	'med + 1	low	"low	-	
high'	-	'high	"trigh	'med - 1	∧ NC(e, V, med)	

Display 3 Specification

Display 3 Program

{Test} if V[med] < e then low := med + 1 else if V[med] > e then high := med - 1 else begin index := med; found := true end

Display 3 Specifications of Invoked Programs

Empty

END OF DISPLAY 3

Display 4 Specification

Initialization	external variables: e, V, found, low, high	(from Display I)
$\mathbf{R}_{\mathbf{i}}(\mathbf{y}) = (low^*)$	= 1) \land (high' = n) \land (found' = false) \land NC(e, V)	

***************************************	1 T C 2 J 1 1
Display 4 Program	

{initialization} iow := 1; high := n; found := false;

> Display 4 Specifications of Invoked Programs <u>Empty</u>

> > END OF DISPLAY 4

25

LEXICON

A. Pascal external definitions and declarations

const n = n; {integer is to be written here} type vector = array[1..n] of integer; var x, j : integer; A : vector; present : Boolean;

INDEX

Name	Used in
А	D0, D1 ₁ , L _A
Body	D1 _{2,3} , D2 _{1,2}
c	D1 _{2,3} , D2 _{1,3} , D3, D4 ₁
Find	D1 _{1,2}
found	D1 _{2.3} , D2, D3, D4
high	D1 _{2,3} , D2, D3, D4
index	D1 _{2.3} , D2 _{1.3} , D3, D4
Initialization	D1 _{2,3} , D4
j	D0, D1 ₁ , L _A
low	D1 _{2,3} , D2, D3, D4
međ	D2 _{2,3} , D3
ħ.	D0, D1 _{1,3} , D4, L _A
present	D0, D1, LA
Test	D2 _{2,3} , D3
۷	DI ₂₃ , D2 ₁₃ , D3, D4 ₁
vector	D0, D1 ₂ , L _A
x	D0, D11, LA

Legend:

• D0

- Di, =1,2,...
- Di_{j} , $i=1,2,..., f \in \{1,2,3\}$
- Di_{jk} , *i*=1,2,..., *j,k* \in {1,2,3}
- L_x, z=A,B,...

denotes Display *i*, denotes Display *i*, part P*j*, denotes Display *i*, parts P*j* and P*k*, denotes the lexicon, part *x*.

denotes the introduction,

October, 1994

Appendix B

"Dutch national flag" example presented on displays

The description of the problem (based on [2], chapter 14), and the discussion of the solution were given in Section 5.2. What follows is the formal documentation for the complete solution. The notation used to specify "buckets" is explained in [8] and [21].

Display 1 Specification

DutchFlag | external variable: v

 $R_0(.) = flag(\vee) \land same_colors('\vee.\vee')$

Display I Program

Procedure declaration:

```
procedure DutchFlag;
var r, w, b : integer;
begin
r := 1; w := N; b := N;
Rearrange(r, w, b)
end {DutchFlag}
```

Display 1 Specifications of Invoked Programs

Rearrange(r, w, b)	external variable: v	(on Display 2)
$R_1(x) = ((r-1) \wedge$		
	⇒	
(partial f	$lag(v', t', w', b') \land (w' = t'-1) \land :same_colors('v, v'))$	ļ

END OF DISPLAY 1

Display 2 Specification

Rearrange(r, w, b) external variable: v

\$_____

(from Display I)

 $\mathbb{R}_1(J) = ((T = 1) \land (W = N) \land (D = N))$

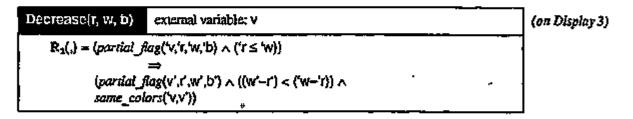
 $(partial_flag(v', r', w', b') \land (w' = r'-1) \land same_colors('v, v'))$

Display 2 Program

Procedure declaration:

procedure Rearrange(var r, w, b : înteger); begin while w ≥ r do Decrease(r, w, b) end {Rearrange}

Display 2 Specifications of Invoked Programs



END OF DISPLAY 2

.........

Display 3 Specification

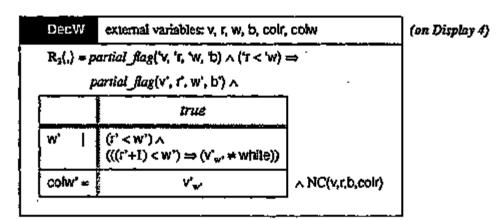
Decrease(r, w, b)	external variable: V	(from Display 2)
R2{,) = {partial_	$lag(V, T, W, D) \land (T \leq W))$	
(partial)	⇒ lag(v',r',w',b') ∧ ((w'-r') < ('w-'r)) ∧	
	ors('v,v'))	· · ·

Display 3 Program

Procedure declaration:

```
procedure Decrease(var r, w, b : integer);
var colr, colw : color;
begin
IncR;
if r < w then begin
DecW;
UseColw
end [if];
UseColr
end [Decrease]
```

Display 3 Specifications of Invoked Programs¹



Display to be continued

¹Note: v, is defined in part C of the lexicon.

IncR	external variables: v, r, w, b, colr	(on Display 5)	•	
$\mathbf{R}_4(\mathbf{x}) = \mathbf{p}$	$cartial_flag(*v, *t, *w, *b) \land (*t \le *w)$	⇒	1 ·	
	partial_flag(v', r', w', b') ∧			
	true		E.	
с 1	$(r' \le w') \land$ $((r' < w') \Rightarrow (v'_{r'} \neq red))$			
coir' =	¥r.	∧ NC(v,w,b,colw)		

UseColr	external va	external variables: V, I, W, D, Colr, Colw				
$R_{s}(.) = par$						
('r :						
	:					
par	tial_flag(v', t', t	$w', b') \land some$	$colors('v,v') \land$			
	ŧ.	'colr =				
	ber	white	eukd	1		
r'	t, = ,t + j	NC(r)	NC(r)			
w' I	NC(w)	w' = 'w − I	W' = W - 1	1		
b' I	NC(b)	NC(b)	b'= b-1	~ NC(coir,colw)		

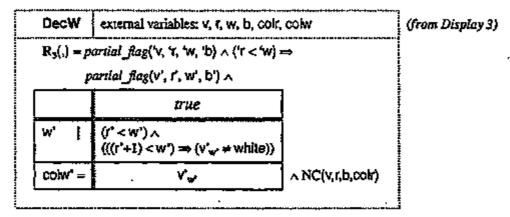
Use	eColw	external va	(on Display 7)			
R ₆ (,			w, "o) ∧ ('colr < 'w) ⇒ ('colv	= 'V.,)	r = 'v-"} ∧	
	parti	al_flag(v', r', v	⇒ w', b') ∧ same_	_calors('V,V'') ∧	(v' _{w'} = colr') ∧	
			'colw =] .	
		red	while	biye		
ŗ	i	f' = 'T + 1	NC(r)	NC(/)	1	
w'	1	NC(w)	w' = 'w - t	w' = 'w - 1	1	
b'		NC(b)	NC(b)	b' = 'b - 1	^ NC(colr,colw)	

END OF DISPLAY 3

c.

1

Display 4 Specification



Display 4 Program

```
{DecW}
colw := LOOK(w);
while (colw = white) and ((r+1) < w) do begin
w := w-1; colw := LOOK(w)
end
```

Display 4 Specifications of Invoked Programs

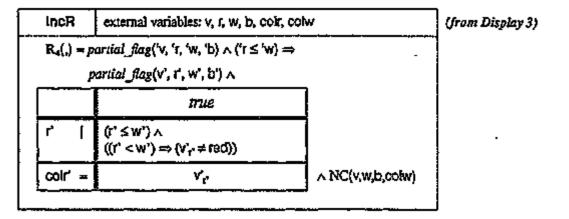
Empty

END OF DISPLAY 4

-)

DISPLAY 5

Display 5 Specification



Display 5 Specification

```
[IncR]

colr := LOOK(r): { v is an implicit variable used by LOOK }

while (colr = red) and (r < w) do begin

r := r+1; colr := LOOK(r)

end
```

Display 5 Specifications of Invoked Programs

Empty

END OF DISPLAY 5

DISPLAY 6

Uş	eCoir	external va	riables: v, r, w,	b, coir, colw		(from Display 3		
R,	$\mathbf{R}_{s}(.) = partial_flag(`v, `r, `w, `b) \land (`colr = `v.,) \land$							
	$('r \leq 'w) \land (('r < 'w) \Rightarrow ('coir \neq red))$							
		-	⇒					
	, par	tial_flag(v', t', v	√', b') ∧ <i>same</i> _	_colors('V,V') ∧	, ,			
		'coir =						
		red	white	blue		• • • •		
r	l	r`='r+1	NC(r)	NC(r)				
w'		NC(w)	$\mathbf{w}^* = \mathbf{w} - 1$	w'='w-1]			
þ.		NC(b)	NÇ(b)	b'='b-1	∧ NC(coir,coiw)	Î T T		

Display 6 Specification

Display 6 Program

REEDINGRAN (REINVERMILTERISTICER

Display 6 Specifications of Invoked Programs

<u>Empty</u>

.....

END OF DISPLAY 6

ï

•

DISPLAY 7

Display 7 Specification

Use	UseColw external variables: v, r, w, b, colr, colw							
R₄(,	$R_{d}(,) = partial_flag('v, 'r, 'w, 'b) \land ('colr = 'v_{\cdot r}) \land ('colw = 'v_{\cdot w}) \land$							
	$(r < w) \land ((r+1) < w) \Rightarrow (colw \neq white))$							
		:	⇒			7		
	partia	il_flag(√, 1', \	N', b') ∧ same_	_colors('V,V') ∧	(V' _{w'} = colr') ^a ∧			
		colw =						
		red	white	blue		-		
					f	1		
r'		$\mathbf{r}' = \mathbf{r} + \mathbf{I}$	NC(r)	NC(r)				
r' w'	 	r' = 'r + 1 $NC(w)$	$\frac{NC(r)}{W' = W - 1}$	NC(r) W' ≠ 'w – 1				

a. The post-condition V'w * Colt' is redundant and has been added for ease of comprehension.

Display 7 Program

```
{UseColw}
case colw of
red: begin SWAP(r, w); r := r+1 end;
white: w := w-1;
blue: begin SWAP(w, b); w := w-1; b := b-1; SWAP(r,w) end
end
```

Display 7 Specifications of Invoked Programs

Empty

END OF DISPLAY 7

LEXICON

A. Auxiliary functions

card: set \rightarrow integer card(s) $\stackrel{\text{df}}{=} |s|$ (i.e. number of elements in the set s)

flag: buckets \rightarrow boolean flag(v) $\stackrel{\text{de}}{=} \exists r, b [partial flag(v, r, r-1, b)]$

 $partial_flag:$ buckets \times integer \times integer \times integer \rightarrow boolean

 $partial_flag(v,r,w,b) \stackrel{\text{df}}{=} (1 \le r) \land (r-1 \le w) \land (w \le b) \land (b \le N) \land$ $\forall i \ (1 \le i \le N) \ [\ ((i < r) \Rightarrow (v_i = \text{red})) \land$ $((w < i \le b) \Rightarrow (v_i = \text{white})) \land$ $((b < i) \Rightarrow (v_i = \text{blue})) \ \}$

Note: v_i is defined in part C of this lexicon.

same_colors: buckets \times buckets \rightarrow boolean

same_colors(v1,v2) ≝

 $\begin{aligned} & (card(\{i \mid (1 \le i \le N) \land (vI_i = red)\}) = card(\{i \mid (1 \le i \le N) \land (v2_i = red)\})) \land \\ & (card(\{i \mid (1 \le i \le N) \land (vI_i = white)\}) = card(\{i \mid (1 \le i \le N) \land (v2_i = white)\})) \land \\ & (card(\{i \mid (1 \le i \le N) \land (vI_i = blue)\}) = card(\{i \mid (1 \le i \le N) \land (v2_i = blue)\})) \end{aligned}$

B. Pascal external definitions and declarations

```
const N = {iteral non-negative integer}
type color = {red, white, blue};
type buckets = {vector(N, color) - cf. part C of this lexicon}
var v : buckets;
procedure LOOK(i : integer);
{cf. part C of this lexicon}
procedure SWAP(i, j : integer);
{cf. part C of this lexicon}
```

C. vector(n,elem) Module Interface Specification

(0) CHARACTERISTICS

- type specified: vector(n,elem)
- · features: single-object, generic
- foreign types: elem, <integer>, <positive_integer>
- generic parameters: n: <positive_integer>, elem

(1) SYNTAX

ACCESS-PROGRAMS

Program Name	Arg#1	Arg#2	Value Type
LOOK	<integer>:V</integer>		elem
PUT	<integer>:V</integer>	elem:V	
SWAP	<integer>:V</integer>	<integer>;V</integer>	

(3) CANONICAL TRACES

canonical(T) \Leftrightarrow T = [PUT(i, e_i)]_{i = 1}^{n} = [PUT(i, ...)]_{i = 1}^{n}

EQUIVALENT NOTATION FOR TRACES

Trace	Equivalent notation
vLOOK(i)>	٧j

(3) EQUIVALENCES

TLOOK(i) => T

T.PUT(i, c) ⇒>

Condition	Equivalence
(1≤i≲n)	%wrong_index%
l≤i≤o	T1.PUT(i,e).T2 where T=T1.PUT(i,x).T2

.

T.SWAP(i, j) ⇒

Condition		Equivalence
$\neg((1 \le i \le n) \land (1 \le j \le n))$		%wrong_index%
	(i < j)	T1.PUT(i,x).T2.PUT(j,y).T3 where T = T1.PUT(i,y).T2.PUT(j,x).T3
$(1 \le i \le n) \land (1 \le j \le n) \land$	(i = j)	T
	(i > j)	T1.PUT(j,x).T2.PUT(i,y).T3 where T = T1.PUT(j,y).T2.PUT(i,x).T3

(4) RETURN VALUES

Program Name	Argument No	Value
LOOK(i)	Valoe	e where vector(n,elem) = T1.PUI(i,e).T2

INDEX

Name	Сатедогу	Used in
b	variable in DutchFlag	Dl ₂
Ъ	formal parameter in Reamange	D1 ₃ , D2 _{1.2}
b	formal parameter in Decrease	D2 ₃ , D3, D4 ₁ , D5 ₁ , D6, D7
blue		D0, D3 ₃ , D6, D7, L _{A,B}
buckets		D0, L _{A.B}
card		L _A
color		D0, D3 ₂ , L _B
·colr		D3 _{2,3} , D4 ₁ , D5, D6, D7 ₁
colw		D3 _{2,3} , D4, D5 ₁ , D6 ₁ , D7
Decrease	•	D2 ₂₃ , D3 _{1,2}
Ds¢W		D3 ₂₃ , D4
DutchFlag		DI ₁₂
flag		Dl ₁ , L _A
IncR		D323, D5
LOOK	L.	D0, D4 ₂ , D5 ₂ , L _{B,C}
N		D0, D123, D21, LAB
partial_flag		D1 ₃ , D2 _{1,3} , D3 _{1,3} , D4 ₁ , D5 ₁ , D6 ₁ , D7 ₁ , L _A
PUT ,		D0, L _c
T	variable in DutchFlag	Dl ₂
r	formal parameter in Reamange	D1 ₃ , D2 _{1,2}
r	formal parameter in Decrease	[•] D2 ₃ , D3, D4, D5, D6, D7
red	· ·	D0, D3 ₃ , D5, D6, D7, L _{AB}
Rearrange		Di ₂₃ . D2 ₁₂
same_colors		D11.3, D213, D313, D61, D71, LA
SWAP		D0, D6 ₂ , D7 ₂ , L _{B,C}
UseCoir		D3 _{2,3} , D6
UseColw		D3 _{2,3} , D7

October, 1994

Name	Category	Used in
v		D0, D1 _{1,3} , D2 _{1,3} , D3 _{1,3} , D4 ₁ , D5 ₁ , D6 ₁ , D7 ₁ , L _B
vector		D0, L _{B,C}
¥	variable in DutchFlag	Dlz
¥	formal parameter in Reamange	D1 ₃ , D2 _{1,2}
¥	formal parameter in Decrease	D2 ₃ , D3, D4 _{1,2} , D5 _{1,2} , D6 _{1,2} , D7 _{1,2}
white		D0, D3 ₃ , D4, D6, D7, L _{A,B}

Legend:

• D0		denotes the introduction,
• Dí, <i>i</i> =1,2,		denotes Display i,
• Di _j , <i>i</i> =1,2,,	$j \in \{1,2,3\}$	denotes Display i, part Pj,
• Dl _{ik} , <i>i</i> =1,2,,	$j,k \in \{1,2,3\}$	denotes Display i, parts Pj and
• L _x , x=A,B,		denotes the lexicon, part x.

d Pk. denotes the lexicon, part x.

October, 1994

-

CRL Report No. 295



Session W1: Natural Language Modeling

Dr. John Sharp Sandia National Laboratories

Natural Language Modeling John K. Sharp, PhD Sandia National Laboratories

This seminar describes a process and methodology that uses structured natural language to enable the construction of precise information requirements directly from users, experts, and mangers. The main focus of this natural language approach is to create the precise information requirements and to do it in such a way that the business and technical experts are fully accountable for the results. These requirements can then be implemented using appropriate tools and technology. This requirement set is also a universal learning tool because it has all of the knowledge that is needed to understand a particular process (e.g., expense vouchers, project management, budget reviews, tax laws, machine function)

Personal accountability for results is established with the expert that is specifying the design and the implementor is accountable for meeting the design requirements. This is done through a systematic procedure based on a common understanding of the requirements and the ability to communicate effectively. In other words, if the craftsman produced the part according to the requirements then he did the correct job. The accountability for form, fit, and function resides with the engineer who created the design. The craftsman is only accountable for meeting the requirements. The center of this accountability process is a communication channel that is completely understood by all of the participants. Natural language modeling processes allow information technology to achieve this same high quality level.

The advantage of this procedure is that it takes an informal, possibly incomplete, possibly redundant, possibly inconsistent and possibly indeterminate description of a user problem and turns it into a precise set of facts and constraints that contain all of the knowledge and business rules that are necessary for completely solving a user problem. The sentences are created and analyzed by the subject matter expert with the analyst being a facilitator or scribe of the knowledge that is created. The expert is fully accountable for the specification and the knowledge can be transformed into desired graphical and textual presentations that become part of the design specification for the implementor.

This seminar will be an overview of the procedure for creating natural language models. Examples will be provided for every step in the procedure. The procedure starts with the subject matter expert verbalizing sentences about the subject area. Placeholders or variables are then assigned within the created sentences. The sentences are then qualified by assigning names to the placeholder and the object. Constraints are then identified and tested. Finally, the results can then be specified in a number of ways, including relational tables. The focus of the seminar shows how low quality initial inputs are turned into high quality requirements that can hold the subject matter expert accountable for the requirements and the implementor accountable for meeting them.

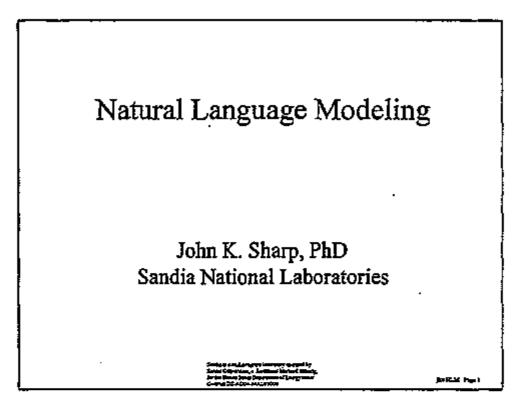
Simple examples will be used throughout the seminar to show how unary, binary and n-ary sentences are analyzed. All possible procedure steps will be presented using examples. Several examples will be used as interactive problems to help attendees understand the procedure.

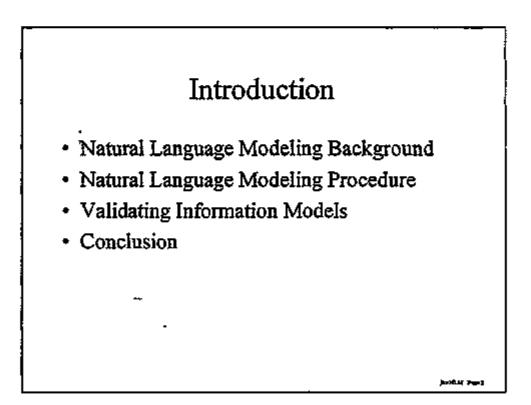
BIOGRAPHY

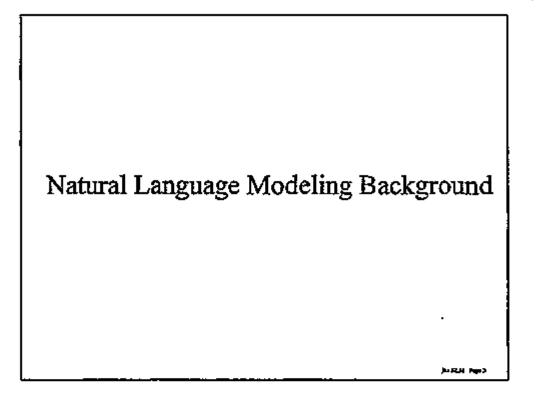
John K. Sharp, PhD

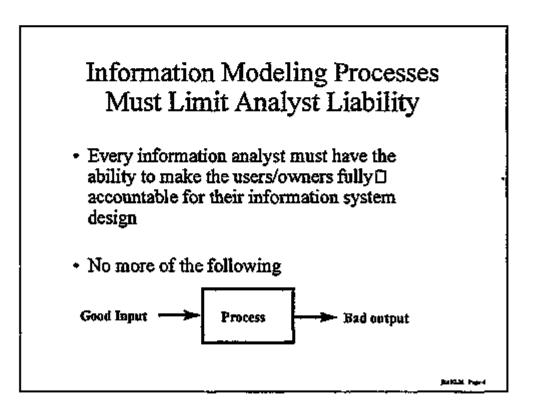
John has performed information analysis in various positions at Sandia for fifteen years. He has worked closely with Prof. Shir Nijssen of the Netherlands for several years to establish the procedure to develop and analyze information problems using structured natural language. They are currently finishing a text on this topic. This procedure was originally based on the NIAM (Natural language Information Analysis Methodology) modeling technique. John and Prof. Nijssen have co-chaired two international conferences on natural language modeling. John is also the editor of the international standard on conceptual schemas.

Sandia National Laboratories Reengineering Center P.O. Box 5800, MS-0803 Albuquerque, NM 87185-0803 Voice: 505-844-5428 Fax: 505-844-7501 E-mail: jksharp@sandia.gov









Natural Language Modeling Overview

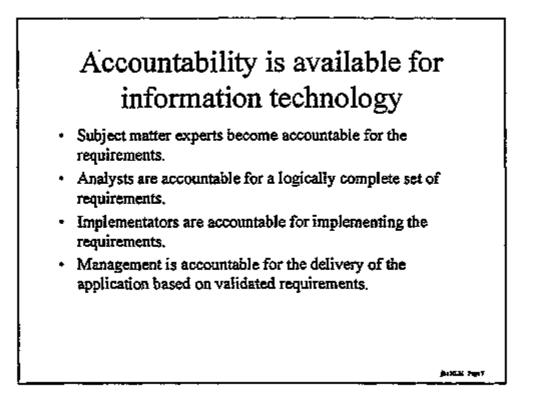
- Based on mathematical analysis of elementary sentences
- Separates analysis from the documentation of analysis
 - Specified analysis procedure that is understandable
 - Can be documented in various graphical models
- Creates a complete design that is validated by subject matter experts
- Accountability can be assigned at every step in the design life-cycle
- Opportunity for significant productivity improvements

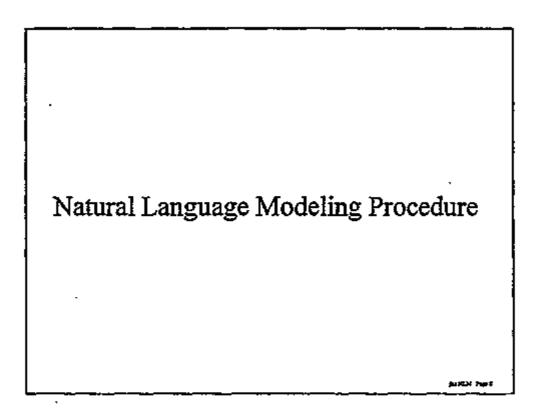
Natural Language Modeling Axioms

- Axiom 1: All the information communicated to and from an information system can be considered to be a set of natural language sentences.
- Axiom 2: In discussions with the user the only language to be used is the familiar jargon of the user.
- Axiom 3: Decisions may only be taken when they are based on a representative number of concrete examples.
- Axiom 4: For every information activity there must be a precise prescription available.

ballini Taged

Auxilia Press





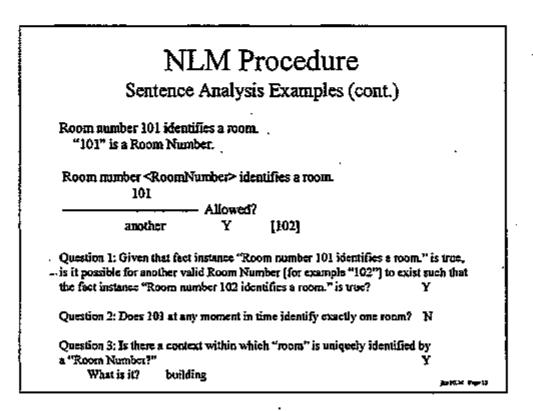
Natural Language Modeling Procedure

- · Sentence analysis questions
- · Sentence analysis examples
- · Sentence analysis procedure
- Process analysis questions
- Process analysis procedure

NLLM Procedure Sentence Analysis Questions Question 1 (Repeated for each variable in a sentence) Given that fact instance "Text a, text." is true, is it allowed for another valid Anr [for example "a,"] to exist such that the fact instance "Text a, text." is true? Question 2 Does a, at *any* moment in time identify exactly one A. Question 3 Is there a context within which A is uniquely identified by an Anr. Question 4 Is there an instance of an identifying fact type that when combined with a, establishes a complete elementary sentence. Where A is an entity or object, Anr is the label, and a, is a population instance.

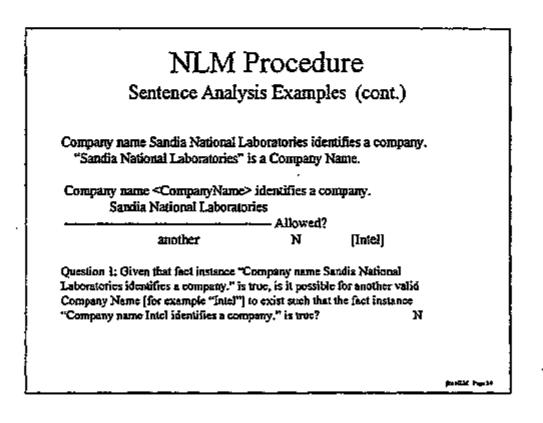
ALTER THE

NLM Pr	oced	ure
Sentence Analy	ysis Ex	amples
Social security number 123-45-67 *123-45-6789" is a Social Secur		-
Social security number <ssn> ide 123-45-6789</ssn>	entaties a p	CESCIL
. 125-40-0787	Allowed?	•
another	Y	[987-65-4321]
Question 1: Given that fact instance "		
identifies a person." is true, is it possil		
Number [for example "987-65-4321"]		
"Social security number 987-65-4321	toenuites e	i person." is true? Y
Question 2: Does 123-45-6789 at any	moment in	time identify exactly one
person?		Y



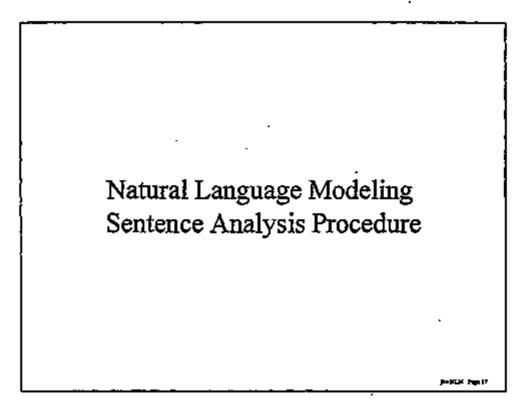
б

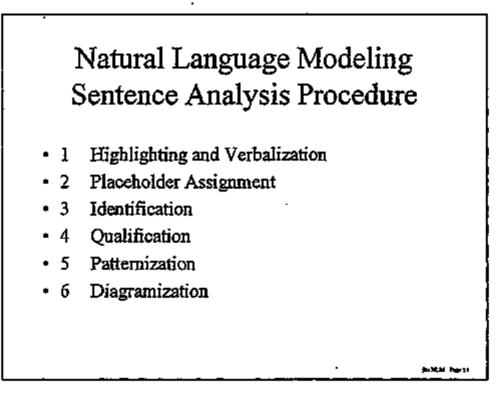
NLM Procedure	
Sentence Analysis Examples (cont.)	
Person name John Smith identifies a person. "John Smith" is an Person Name.	
Person name <person name=""> identifies a person. John Smith Allowed?</person>	
another Y [Suc Jones]	
Question 1: Given that fact instance "Person name John Smith identifies a person." is it possible for another valid Person Name [for example "Sue Jones"] to exist sur fact instance "Person name Sue Jones identifies a person." is true?	-
Question 2: Does John Smith at any moment in time identify exactly one person?	N
Question 3: Is there a context within which "person" is uniquely identified by a "F Name?"	erson N
Question 4: Is there an instance of an identifying fact type that when combined wi name establishes a complete elementary sentence? What is it? Social security number 123-45-6789 identifies a person.	th person Y
	AND THE O



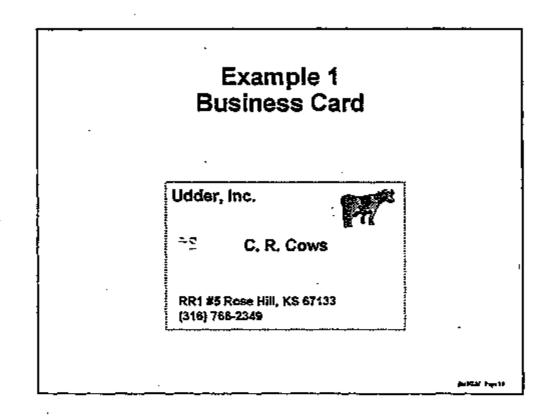
DECOMPOSED Sentence Analysis Examples (cont.) The preceding examples were all unary sentences (only one placeholder in the sentence can vary). The only additional requirement for binary or higher order sentences is to extend the first question to allow each placeholder to independently vary. This is done by creating a matrix of the valid instance and replacing the instance values on the diagonal with "another." Question 1 is then asked for each of the sentences. Questions 2 - 4 are asked exactly like they were in unary sentences.

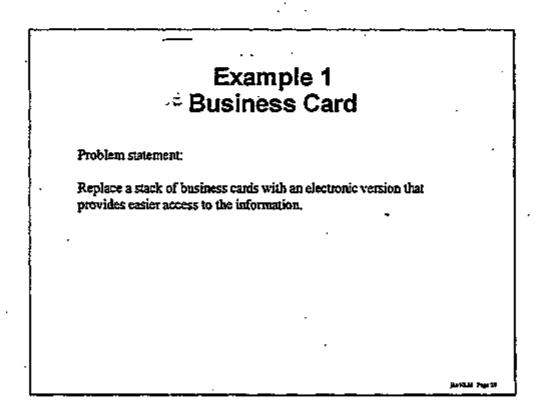
		NL	M Pro	cedur	e	
	Se	ntence A	malysis E	xamples	(cont.)	
Room number 1 "101" is a Ro "803" is a Bu Room number - 10	om i ildir Ro	Number. Ig Id.			· identifies a ro	оп.
			Allowed?			
anoti 10		803 another	Y Y	[102] [801]	-,	
Question 3.1: Giv is true, is it possib						
the fact instance "					-	Ŷ
Question 1.2: Giv is true, is it possib	le fo	r another vali	d Building Id (1	fot example "	80)"] to exist su	ch that
the fact instance "			-			Y
Question 2: Does	101	in 803 et any	moment in time	identify exa	ctly one room?	Y
						Jan Ki, M. Page 34

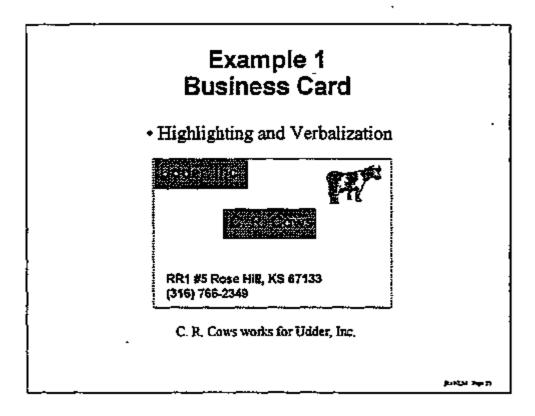


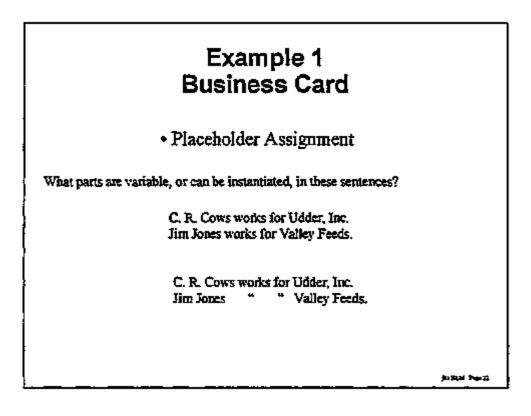


....



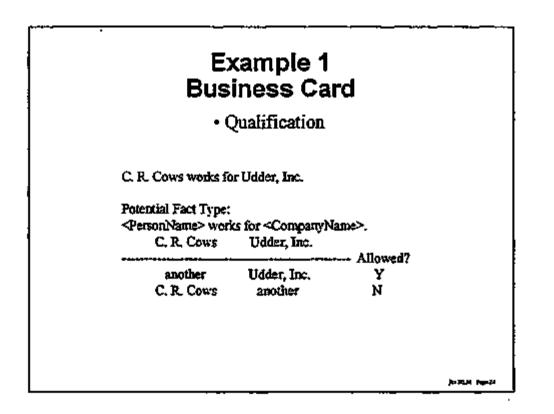


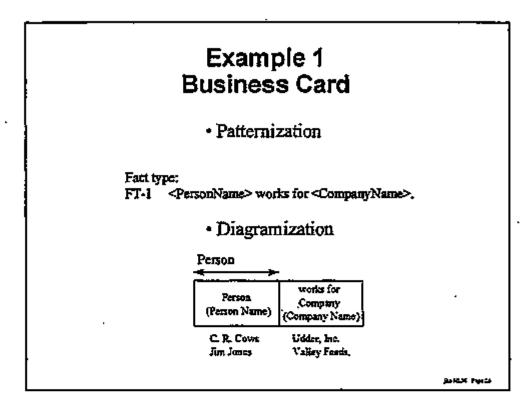




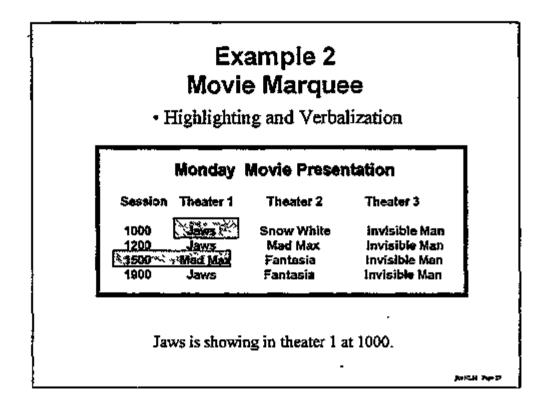
 \mathbf{n}

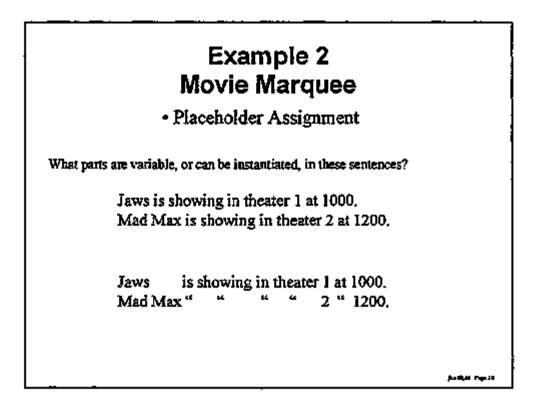
Exampl Business		
 Identificat 	ion	
C. R. Cows works for U Jim Jones " " V	ldder, Inc. /alley Feeds.	
Of which class are C. R. Cows and Jim Jon Of which class are Udder, Inc. and Valley H		Person Company
How is an individual element of the population Person Name How is an individual element of the population Company Name	•	
What is the name of the placeholder for the j Jim Jones appear in this sentence? What is the name of the placeholder for the j	<pre><personname< pre=""></personname<></pre>	>
Valley Feeds appear in this sentence?	<company:na< td=""><td>me></td></company:na<>	me>





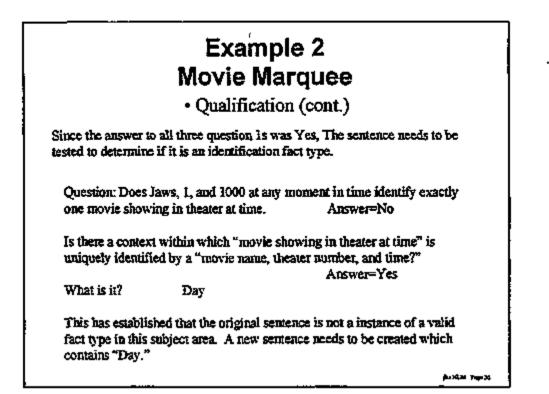
		ie Marque	
	Monday	Movie Preser	ntation
Session	Theater 1	Theater 2	Theater 3
1000	Jaws	Snow White	invisible Mar
1200 1500	Jaws Mad Max	Mad Max Fantasia	Invisible Mar Invisible Mar
1900	Jaws	Fantasia	Invisible Mar

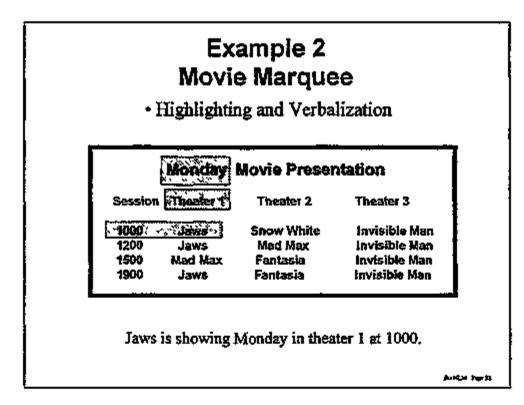




Example 2 Movie Marque	e
Identification	
Jaws is showing in theater 1 at Mad Max " " " 2 "	1000. 1200.
Of which class are <> elements? Jaws and Mad Max l and 2 - 1000 and 1200	Movie Theater Time
How is an individual element of the population of the cla	uss <> identified?
Movie Movie Name Theater Theater Number Time Time	
What is the name of the placeholder for the position whe	re <> appear in this sentence?
Jaws and Mad Max MovieNa I and 2 TheaterN 1000 and 1200 Time	

	Movie	ample 2 e Marqu		
	• Qu	alification		
Jaws is showing in t	heater 1 at 1000).		
Potential Fact Type: <moviename> is si Jaws</moviename>	-	1000	iber> at <tin Allowed?</tin 	ae>.
another	1	1000	Y	
Jaws	another	1000	Ŷ	
Jaws	1	another	Y	
Question: Given tha is it allowed for ano the fact instance "M	ther valid Movie	e [for example]	"Mad Max"']	-
				Jan Kilal, Page 30

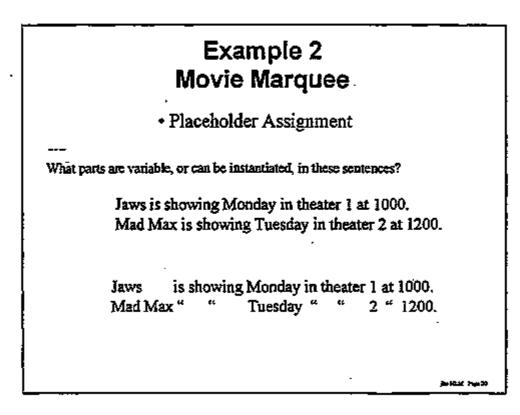




I

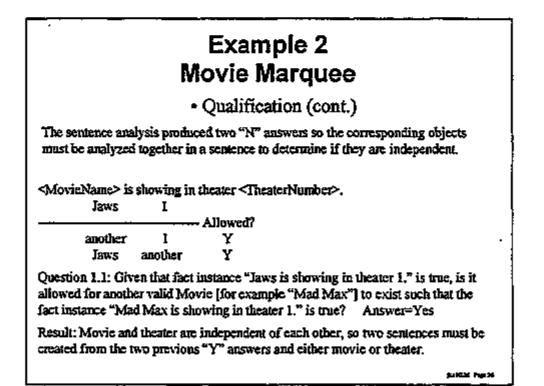
1

ì



Example 2 Movie Marquee
 Identification
The new variable must now be identified.
Jaws is showing Monday in theater 1 at 1000. Mad Max " " Tuesday " " 2 " 1200.
Of which class are Monday and Tuesday elements? Day
How is an individual element of the population of the class Day identified? Day
What is the name of the placeholder for the position where Monday and Tuesday appear in this sentence? Day
ja 2022 7mg 34

		Examp vie Ma	rquee	
	•	Qualifica	tion	
laws is showing Me	onday in theat	er I at 1000.		
Potential Fact Type	:			
MovieName> is s		> in theater <	TheaterNum	ber> at <time>.</time>
Jaws	Monday	1	1000	
<u>.</u>				Allowed?
another	Monday	1	1000	N
Jaws	another	1	1000	Y
ZwsĽ	Monday	another	1000	N
Jaws	Monday	1	another	Y
s true, is it allowed	for another va	aliđ Movie (f	or example "	onday in theater 1 at 1000 'Mad Max"] to exist such ater 1 at 1000." is true? 0



Movie Marquee • Qualification (cont.) Jaws is showing Monday at 1000. Potential Fact Type: Morday 1000 Allowed? another Monday 1000 Jaws Monday 1000 another Monday 1000 Y Jaws another 1000 Y Jaws another 1000 Y Jaws Monday another Y Question 1.1: Given that fact instance "Jaws is showing Monday at 1000." is true, is it allowed for another valid Movie [for example "Mad Max"] to exist such that the fact instance "Mad Max is showing Monday at 1000." is true? Answer=Yes Question 2: Does Jaws, Monday, and 1000 at any moment in time identify exactly one movie showing on day at time. Answer=Yes		E	xampi	le 2	
Jaws is showing Monday at 1000. Potential Fact Type: MovieName> is showing <day> at <time>. Jaws Monday 1000 another Monday 1000 Y Jaws another 1000 Y Jaws Monday another Y Question 1.1: Given that fact instance "Jaws is showing Monday at 1000." is true, is it allowed for another valid Movie [for example "Mad Max"] to exist such that the fact instance "Mad Max is showing Monday at 1000." is true? Answer=Yes Question 2: Does Jaws, Monday, and 1000 at any moment in time identify exactly</time></day>		Μον	vie Ma	rquee	•
Potential Fact Type: < <u>MovieName></u> is showing < <u>Day></u> at < <u>Time></u> . Jaws Monday 1000 <u>another</u> Monday 1000 Y Jaws another 1000 Y Jaws Monday another Y Question 1.1: Given that fact instance "Jaws is showing Monday at 1000." is true, is it allowed for another valid Movie [for example "Mad Max"] to exist such that the fact instance "Mad Max is showing Monday at 1000." is true? Answer=Yes Question 2: Does Jaws, Monday, and 1000 at any moment in time identify exactly		• Qu	alificatio	n (cont.)	
<moviename> is showing <day> at <time>. Jaws Monday 1000 another Monday 1000 Y Jaws another 1000 Y Jaws Monday another Y Question 1.1: Given that fact instance "Jaws is showing Monday at 1000." is true, is it allowed for another valid Movie [for example "Mad Max"] to exist such that the fact instance "Mad Max is showing Monday at 1000." is true? Answer=Yes Question 2: Does Jaws, Monday, and 1000 at any moment in time identify exactly</time></day></moviename>	Jaws is showing Mo	onday at 1000.	•		
another Monday 1000 Y Jaws another 1000 Y Jaws Monday another Y Question 1.1: Given that fact instance "Jaws is showing Monday at 1000." is true, is it allowed for another valid Movie [for example "Mad Max"] to exist such that the fact instance "Mad Max is showing Monday at 1000." is true? Answer=Yes Question 2: Does Jaws, Monday, and 1000 at any moment in time identify exactly	<moviename> is si</moviename>	howing <day></day>			
Jaws Monday another Y Question 1.1: Given that fact instance "Jaws is showing Monday at 1000." is true, is it allowed for another valid Movie [for example "Mad Max"] to exist such that the fact instance "Mad Max is showing Monday at 1000." is true? Answer=Yes Question 2: Does Jaws, Monday, and 1000 at any moment in time identify exactly	another	Monday	1000		
Question 1.1: Given that fact instance "Jaws is showing Monday at 1000." is true, is it allowed for another valid Movie [for example "Mad Max"] to exist such that the fact instance "Mad Max is showing Monday at 1000." is true? Answer=Yes Question 2: Does Jaws, Monday, and 1000 at any moment in time identify exactly	Jaws	another	1000	Y	
it allowed for another valid Movie [for example "Mad Max"] to exist such that the fact instance "Mad Max is showing Monday at 1000," is true? Answer=Yes Question 2: Does Jaws, Monday, and 1000 at any moment in time identify exactly	Jaws	Monday	another	Y	
	it allowed for anoth	er valid Movie Max is showie	e [for examp ng Monday a	le "Mad Max" t 1000," is tri	"] to exist such that the e? Answer=Yes

	м	Exan Iovie I	nple : Marqi		. 0
	•	Qualific	ation (cont.)	
Jaws is showing	•	-	``	*	
Potential Fact Ty Theater < Theater 1		1000	Day> at <1 Allowed?	Time>.	
another	Monday	-	Y		
1	another	1000	Y		
1	Monday	another	Y		
	nother valid	Theater [for	r example	"2"] to exis	fonday at 1000." is true, at such that the fact Answer=Yes
Question 2: Does one theater in us	-		0 at any n	soment in ti	me identify exactly Answer=Yes
l	_				Justial Page Sa

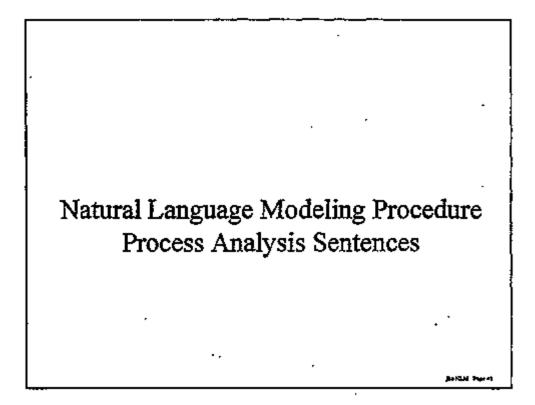
	Exam	ple 2		
	Movie N	larquee		
•	 Pattern 	nization		
Fact Type: FT1 <moviename></moviename>			:rNumber> a	ti <time≻.< th=""></time≻.<>
	•	amization		
Movie_Da	y_Tîme	,		-
	·			Ξ
Movie Name	Day	Theater Number	Time	
Jaws	Monday	1	1000	
Snow W	hite Monday	2	1000	
Mad Ma:	x Tuesday	1	1200	

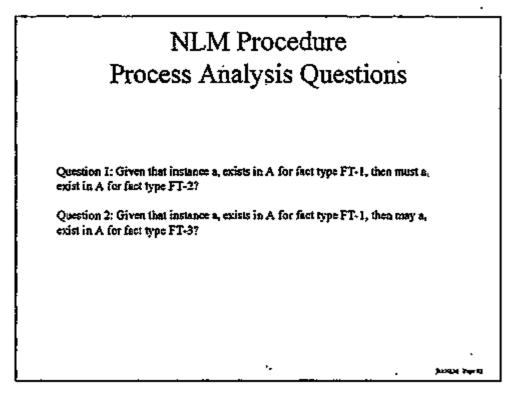
Person Week Case Charge Default Fri Sat SumMon Tue Wed Tinu Approver												
Person	Week Ending	Case	Charge Type	Default Case	Fri	Sat	Sm	Mon	Tuc	Wed	Tinu	Арргочет
1734	1/7/97	23-0	R .		6			- 7			4	5464
1234	1/1/97	662	R		1 2	1		1	8		4	5464
1234	1091	2040	0		5	1		— ,				
1204	10497	23-0	R		1			3	Σ		8	6754
5464	1797	2941	R					8		1	4	5464
1234	1/21/97	Z341	ĸ	X	1			I	8	8	8	5464
1342	Intert	431	F			8				1		5464
1343	เกต	2341	ĸ					8		1 *		3464
2144	17197 .	2451	R		4			4	8	4		5464
6754	inat	4021	R		I					8	1	6534
1342	17) #57	5461	R	×	r—		i —		8	*		5464
13401	1/1497	540	R				Ī					5444
કામ 🗌	1787	201	R		4	i		1	4	4	4	60

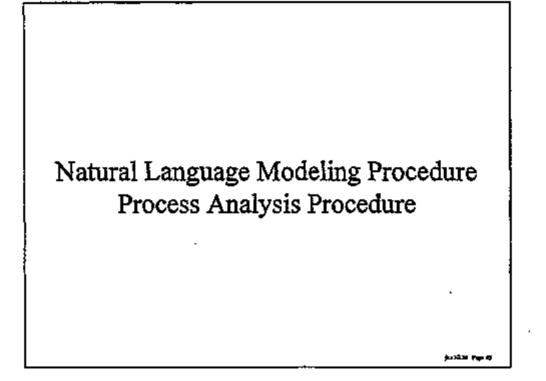
₽

20

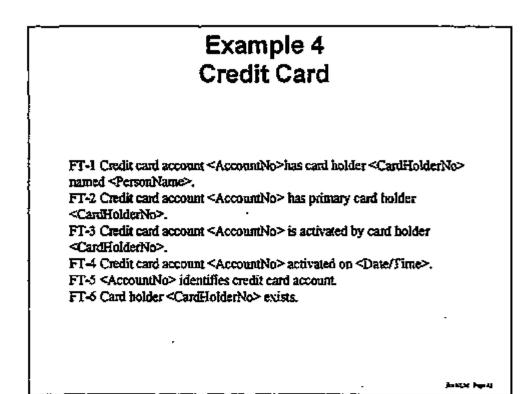
^)

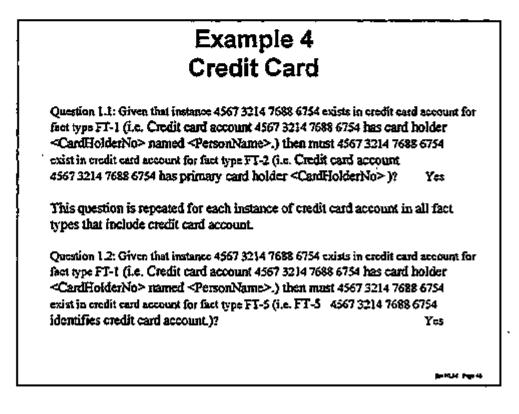


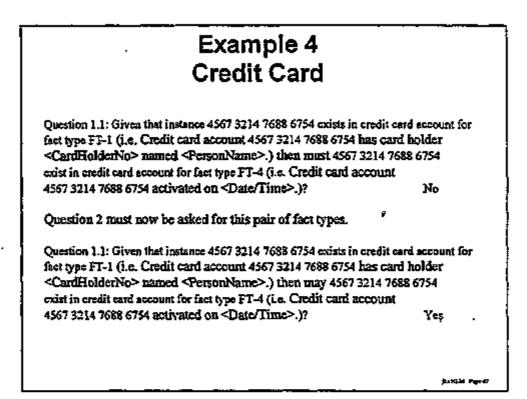


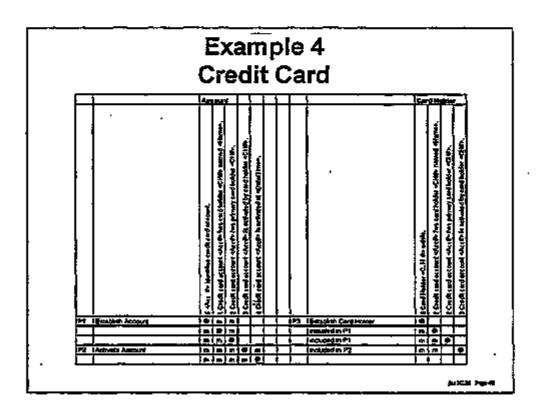


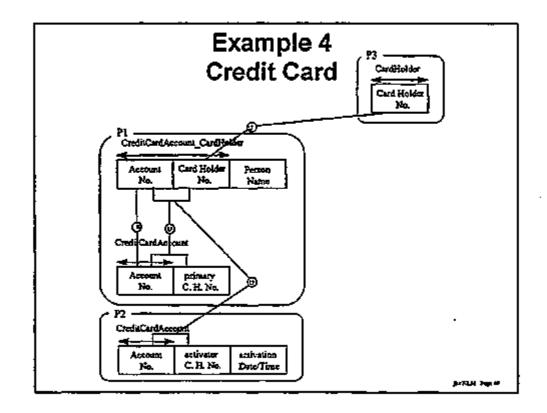
NLM Procedure Process Analysis Procedure • 1 Mandatory • 2 Exclusion

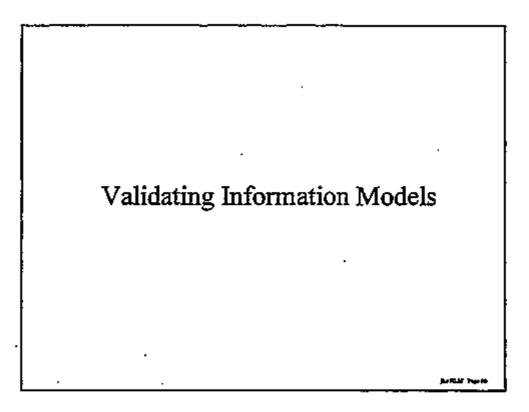


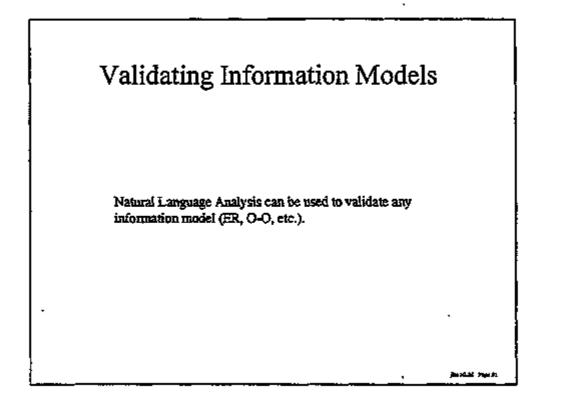


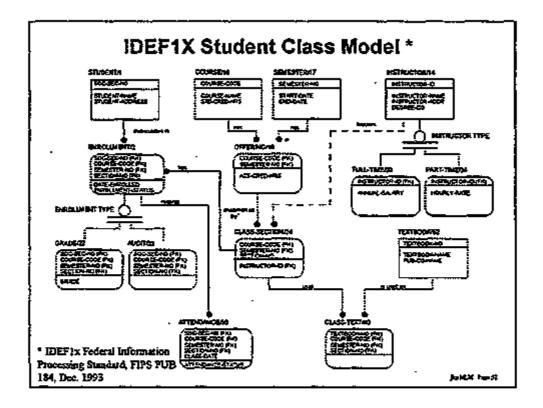


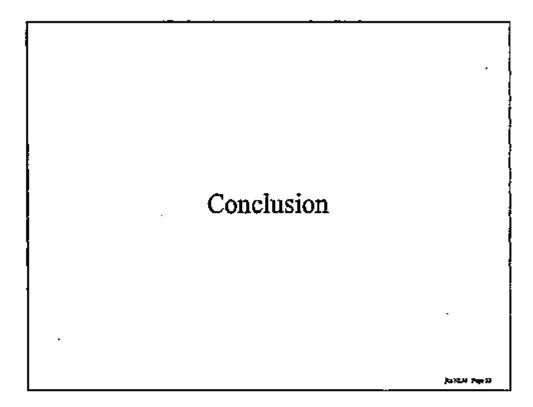


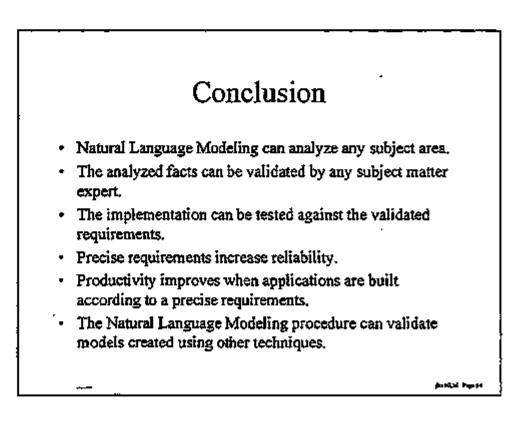


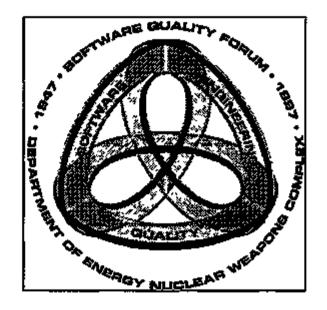












Session X1: Defining Software Processes

Dr. Gerald McDonald Consultant, Sandia National Laboratories

DEFINITION AND DOCUMENTATION OF ENGINEERING PROCESSES

GERALD W. MCDONALD, Ph.D.

This tutorial is an extract of a two-day workshop developed under the auspices of the Quality Engineering Department at Sandia National Laboratories. The presentation starts with basic definitions and addresses why processes should be defined and documented. It covers three primary topics: (1) process considerations and rationale, (2) approach to defining and documenting engineering processes, and (3) an IDEF0 model of the process for defining engineering processes.

Process considerations and rationale introduce models for documenting processes; describe the general architecture for product development; and define implications of immature processes versus those for mature processes.

The approach describes the top-level subprocesses that make up the methodology for definition and documentation of engineering processes; namely: planning, gaining management approval for a process definition project, collecting data on the as-is process to capture current best practices within the organization, constructing a model of the as-is process, and verifying and validating that model.

The final portion presents a four-level, hierarchical model that describes HOW to define and document an engineering process.

BIOGRAPHY

GERALD W. MCDONALD, Ph.D.

Dr. McDonald has a Bachelor of Science in Engineering Science and a Master of Science in Computer Systems Management from the Naval Postgraduate School. Following his retirement the Navy he received a Master of Engineering in Industrial and Systems Engineering and a Ph.D. in Quantitative Management Science (Operations Research) from the University of Florida.

Following receipt of his Ph.D. he worked for BDM International as an executive-level Program and/or Project Manager and technical leader. During his thirteen years with that firm he led both software and non-software projects.

During the three years since his retirement from BDM he has acted as consultant to Sandia, SEMATECH, and a number of other organizations. As a consultant he has, worked primarily in the field of Software Process Improvement. Besides direct technical assistance he has presented training and workshops in software areas such as: quality engineering, software inspections, process definition and documentation, and metrics.

9220 Masini Lane, NW Albuquerque, NM 87114-6001 Voice: 505-898-3277 E-mail: GeraldWMcDonald@juno.com

Definition and Documentation of Engineering Processes

(Tutorial)

Gerald W. McDonald, Ph.D. 9220 Masini Lane, NW Albuquerque, NM 87114 (505) 898-3277

State of Software Practice

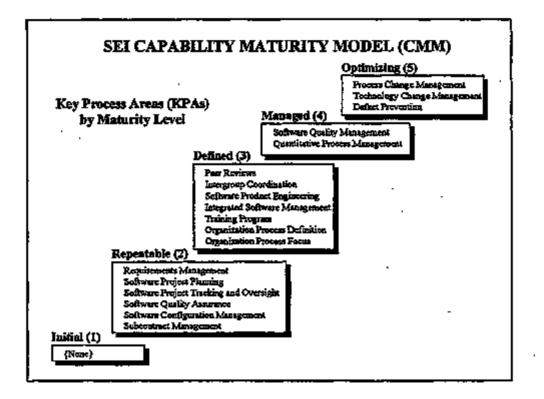
So many software projects fail in some major way that we have <u>had to redefine "success"</u> to keep everyone from being despondent.

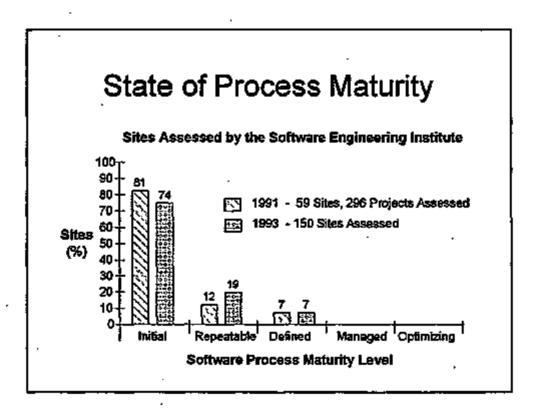
Projects are sometimes considered <u>successful when the</u> <u>overruns are held to 30%</u>, or when the <u>user only junks a</u> <u>guarter</u> of the result.

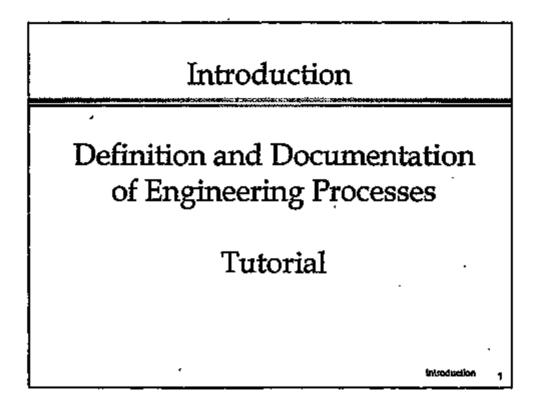
Software personnel are often willing to <u>call such efforts</u> successes.

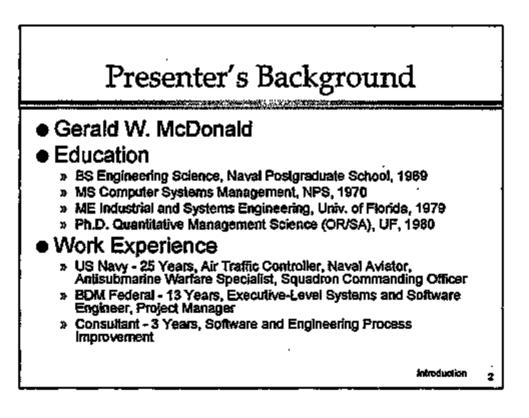
Members of our user community are less forgiving. <u>They know failure when they see it!</u>

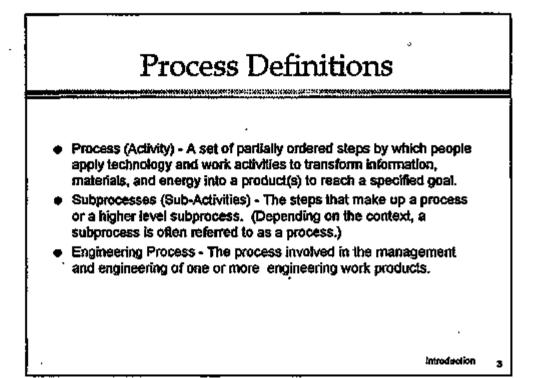
"Controlling Software Projects," by Tom DeMarco

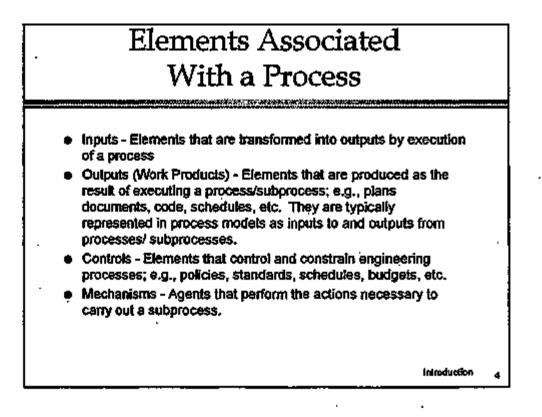


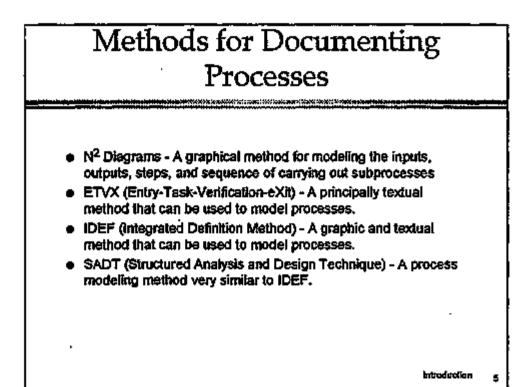


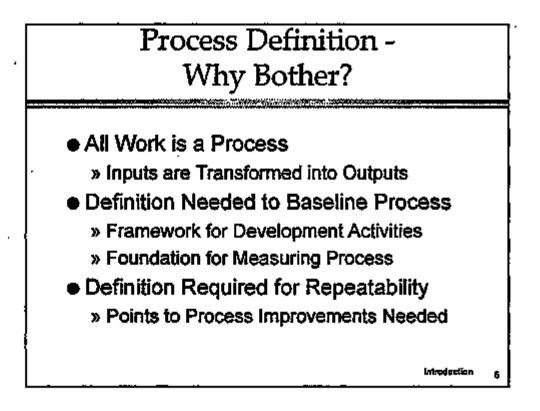


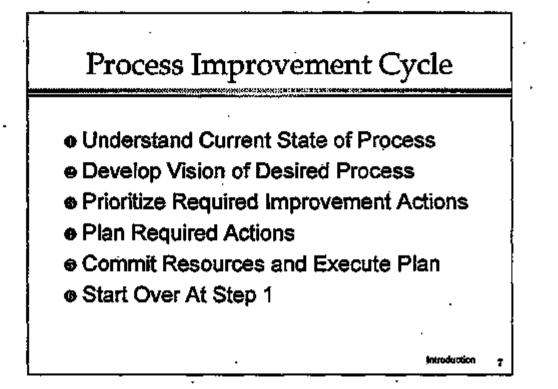


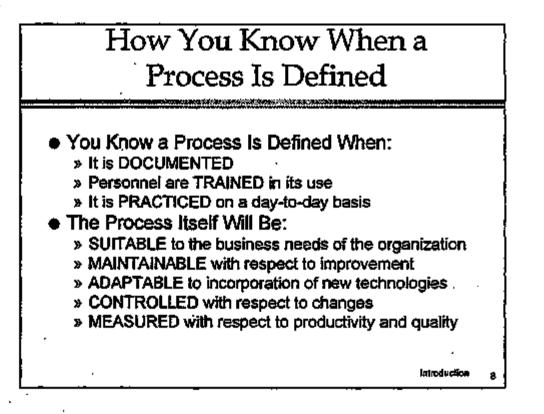


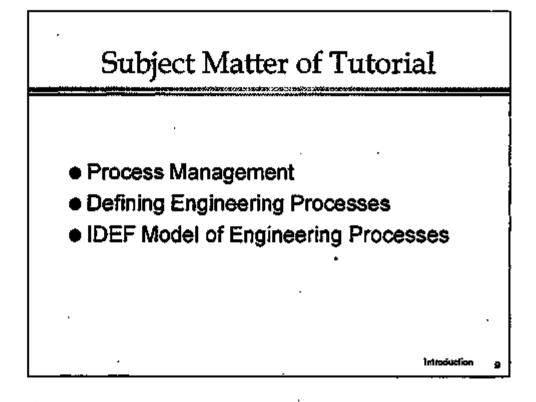


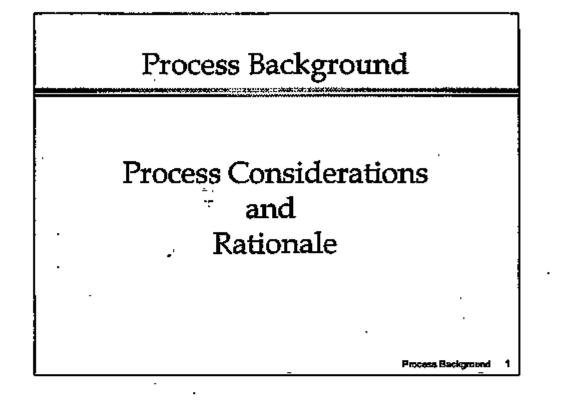


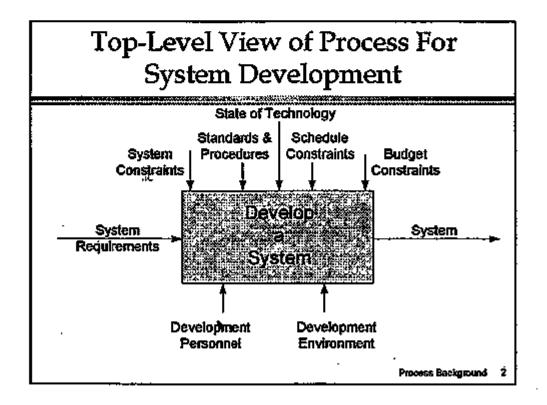


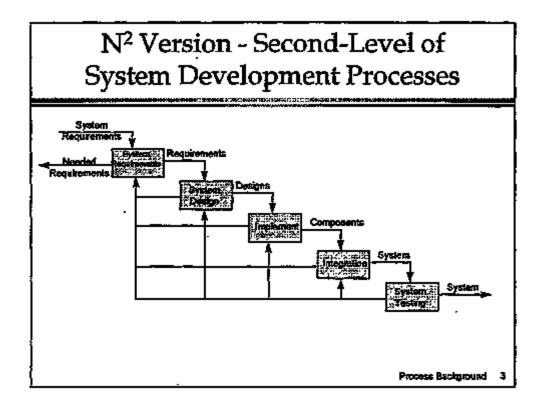


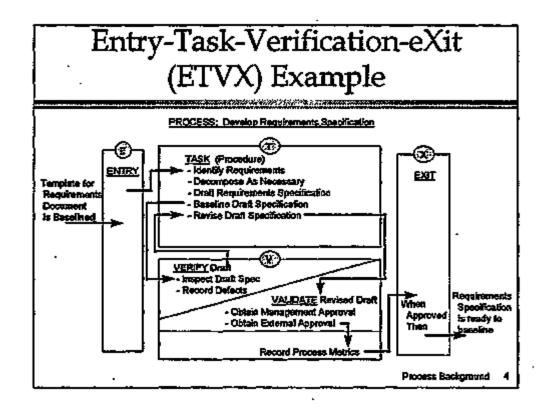


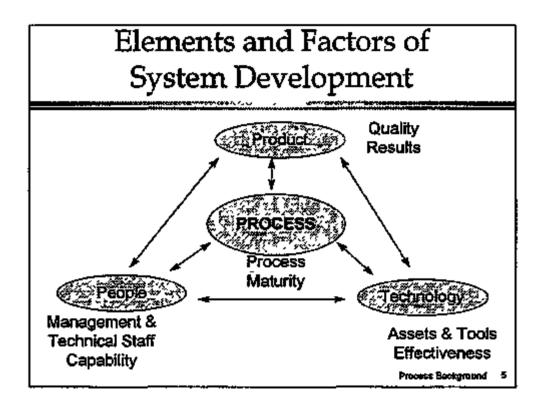


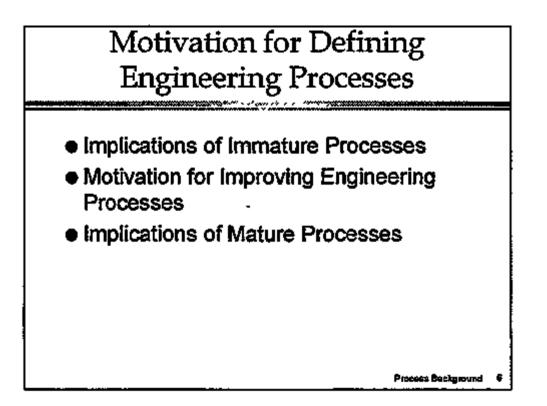


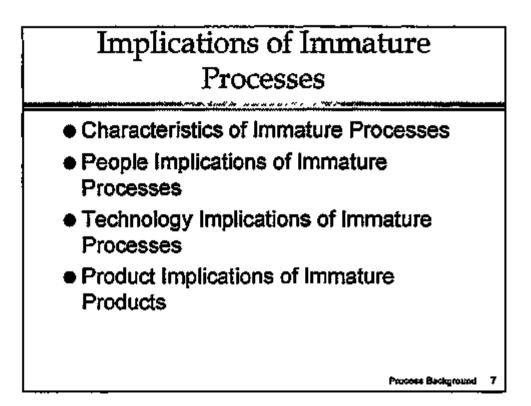


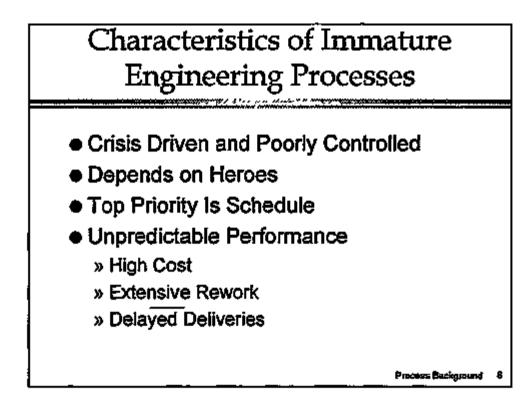


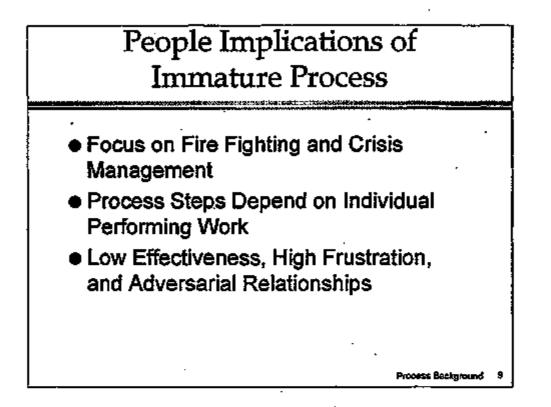


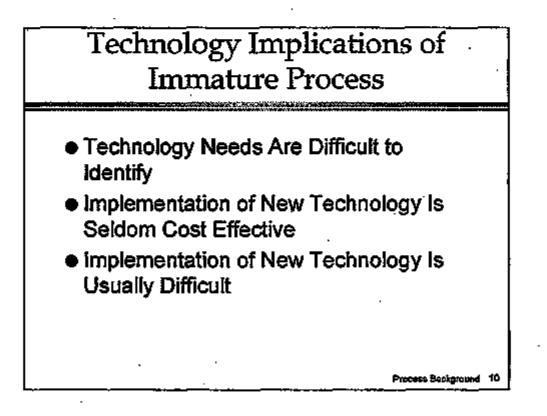


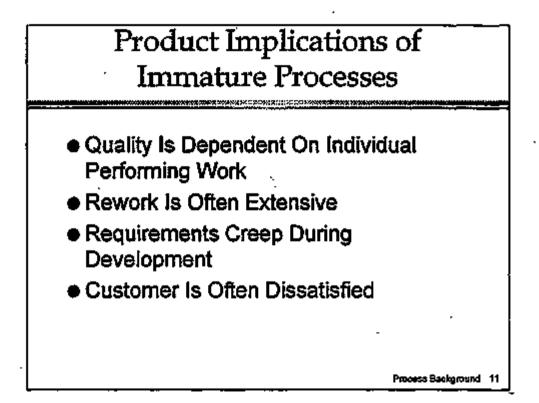


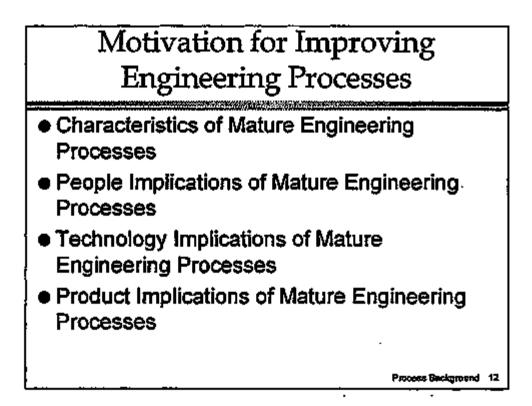


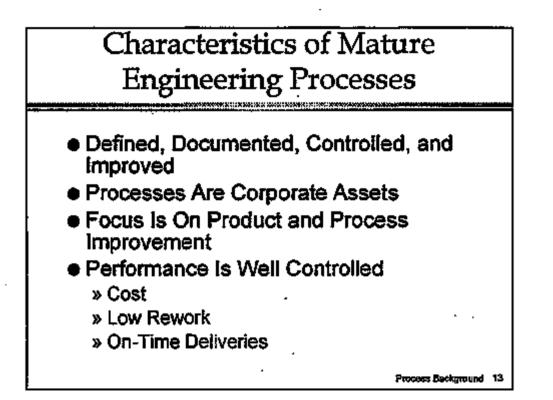


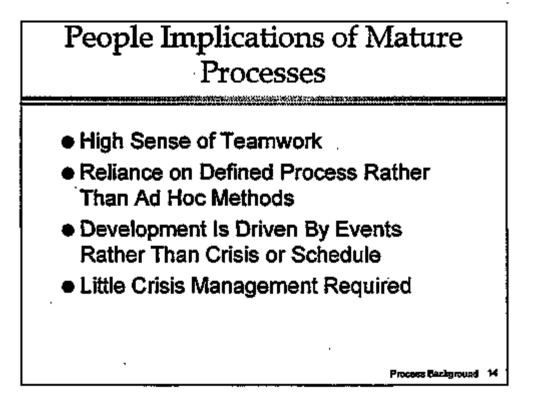












Technology Implications of Mature Processes

- Technology Needs Can Be Identified
- Quantitative Basis Can Be Developed to Support Automation Needs and Selection
- Potential Impacts of New Technology Can Be Estimated More Accurately

Product Implications of Mature Processes

- Cost of Quality Is Very Low and Independent of Individuals Performing Work
- Customer Is More Often Satisfied With Products
- Rework Requirements Are Often Negligible

Process Beckground 16

Process Background 15

Approach to Defining Engineering Processes

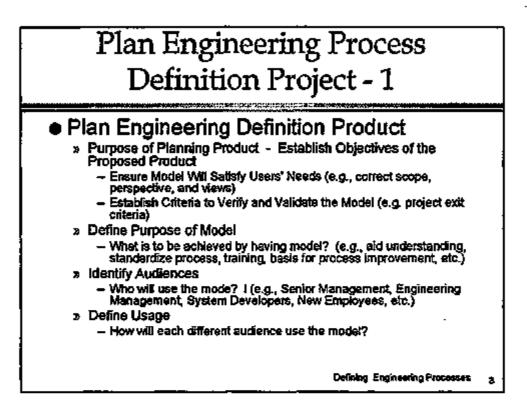
Preparations for and Modeling of Engineering Processes

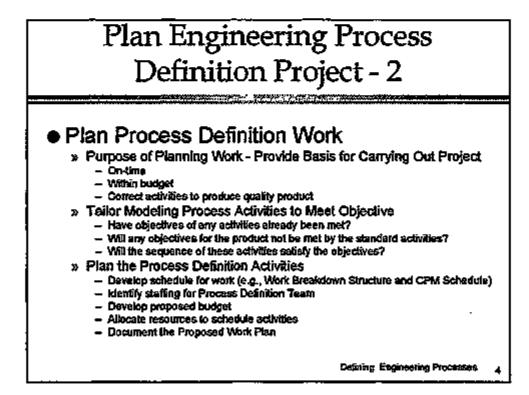
Defining Engineering Processes

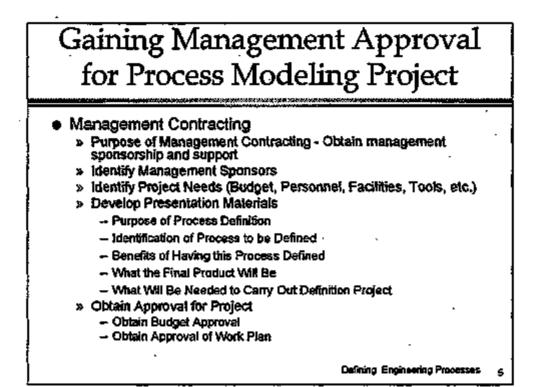
Outline - Approach to Defining Engineering Processes

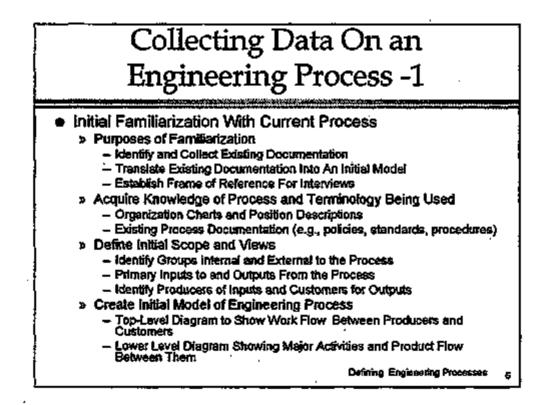
- Prepare for Engineering Process Modeling -
 - » Plan Process Definition Project
 - » Gain Management Approval
- Model Engineering Process
 - » Collect Data on Engineering Process
 - *» Construct Engineering Process Model
 - » Verify and Validate Process Model

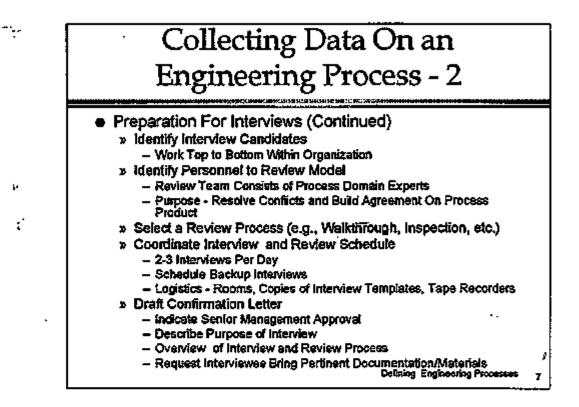
Defining Engineering Processes 2

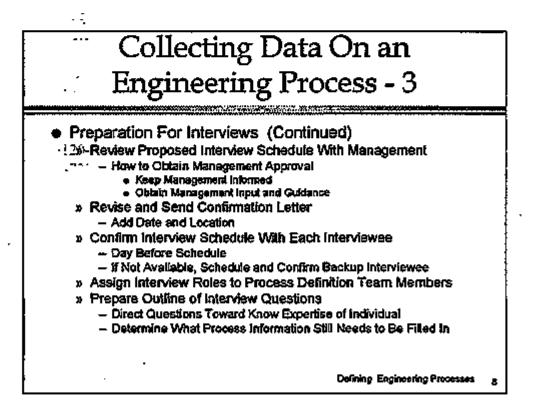


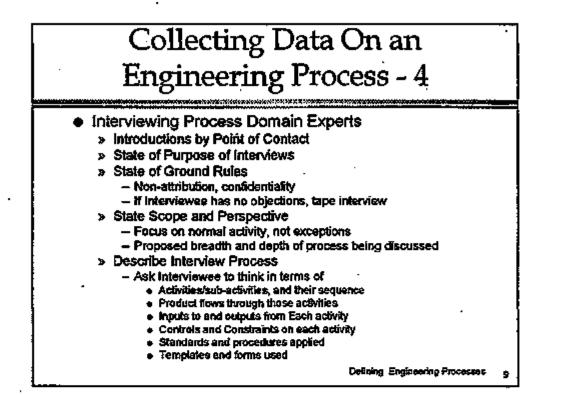


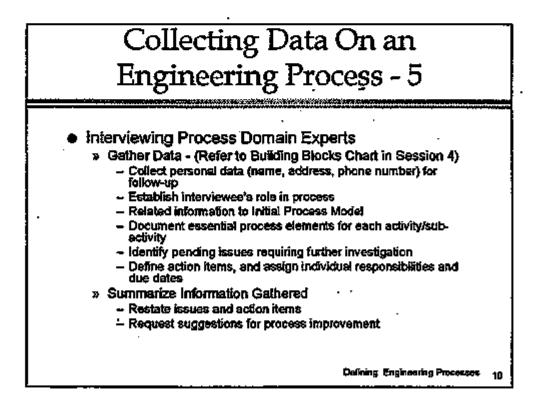


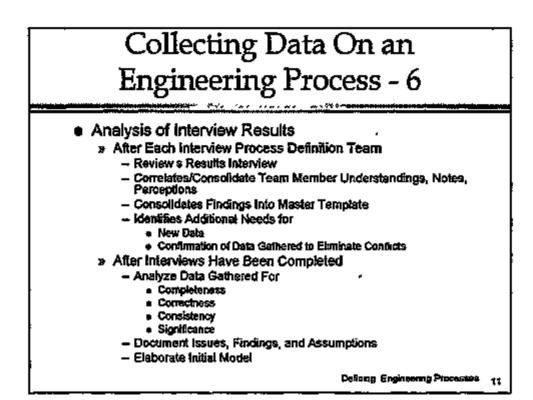


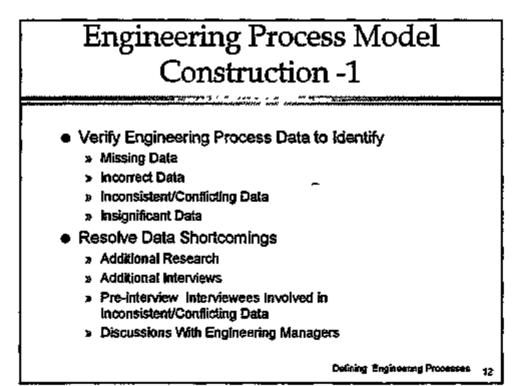


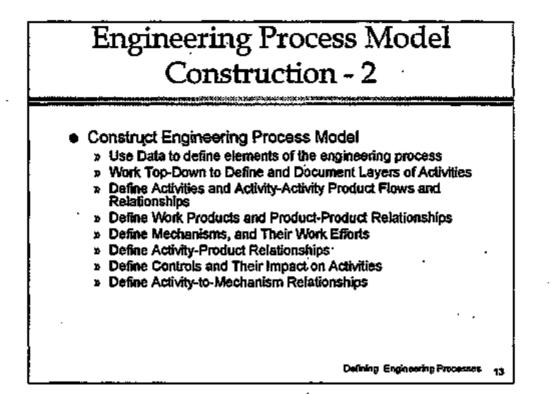


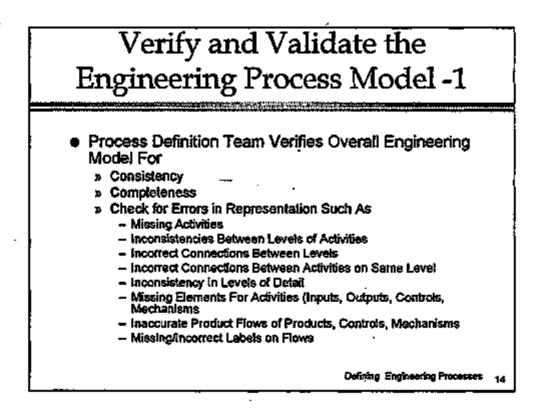


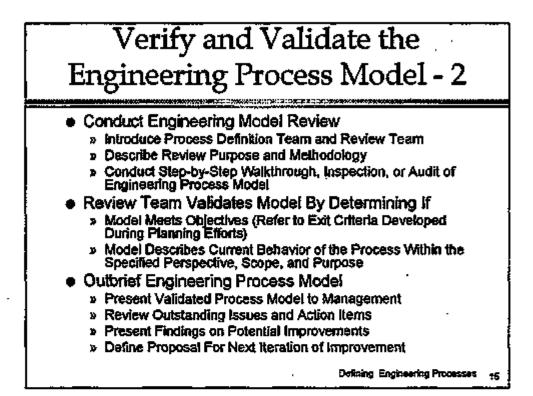


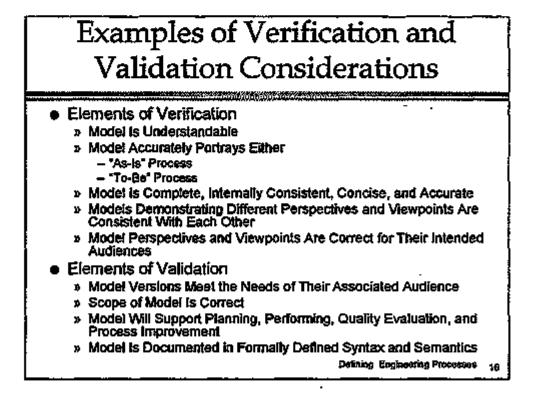


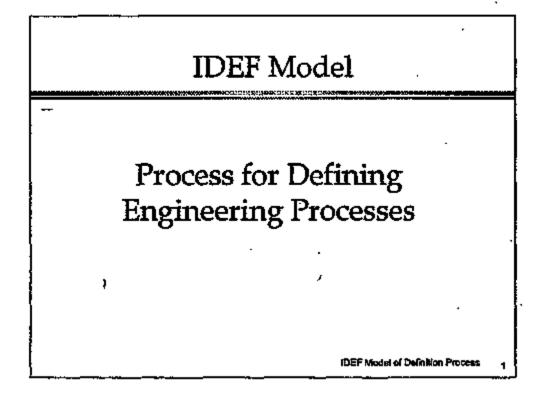


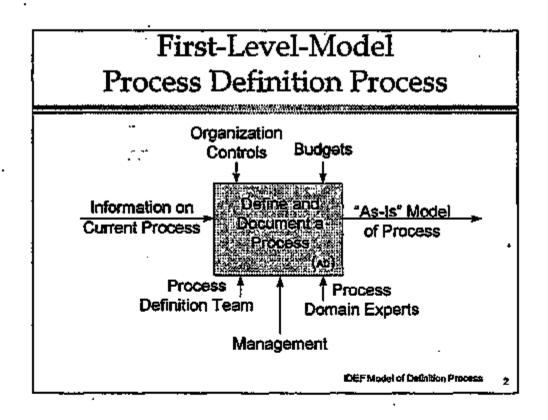


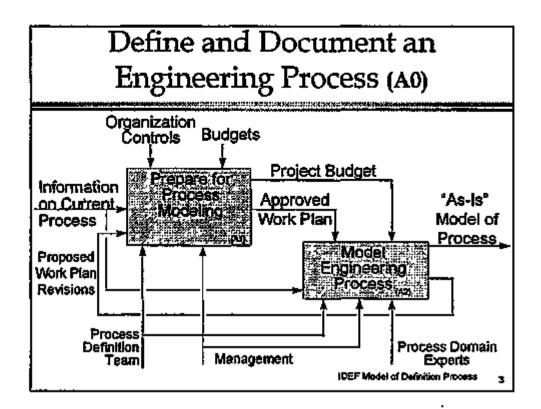


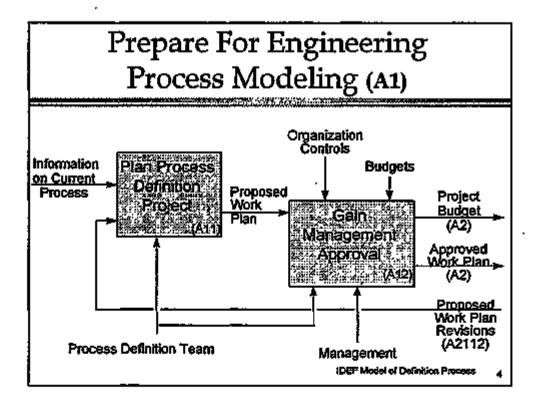


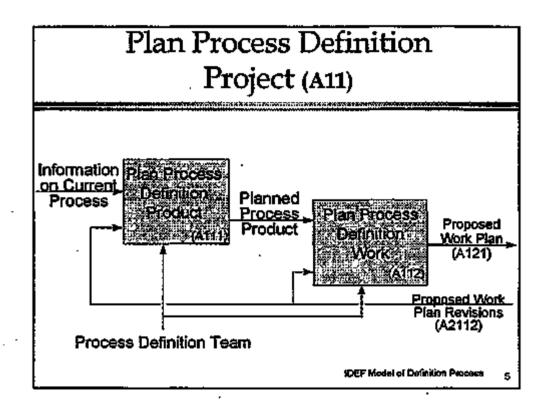


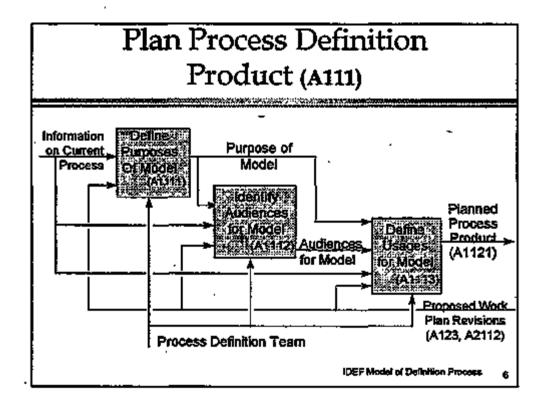


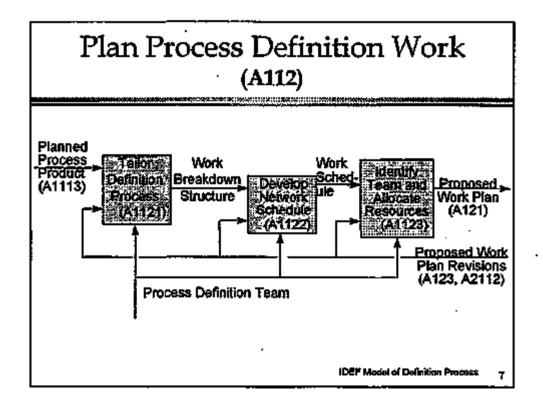


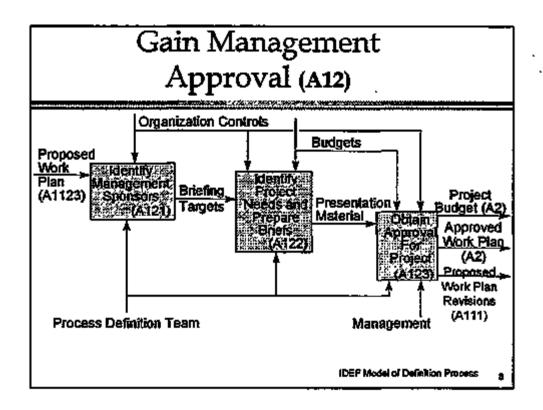


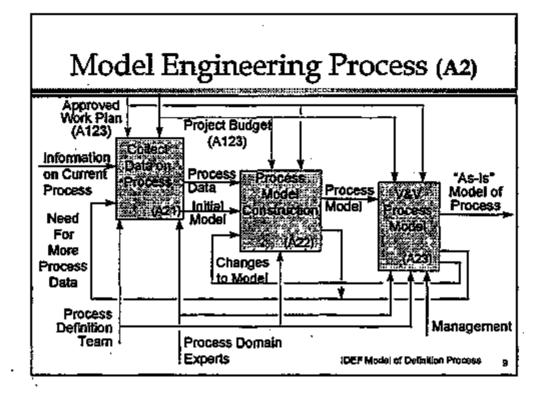


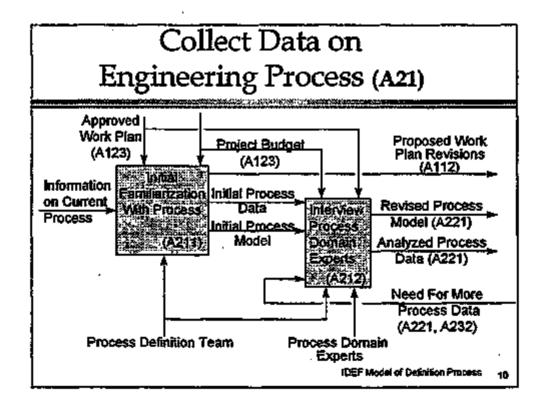


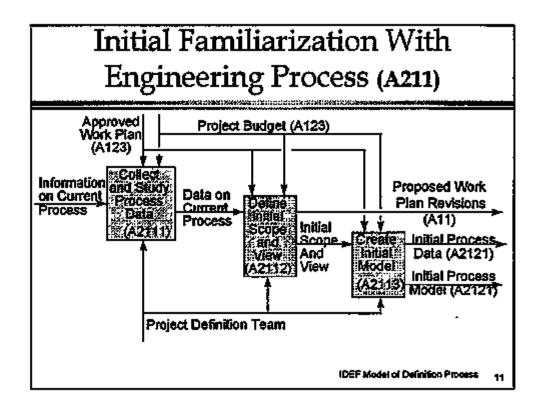


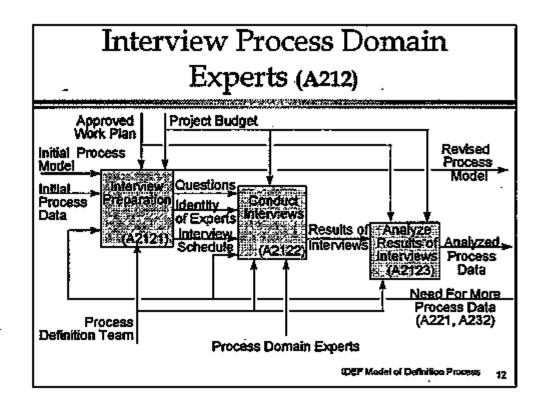


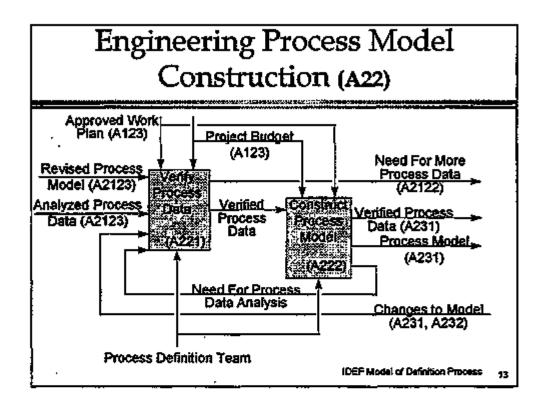


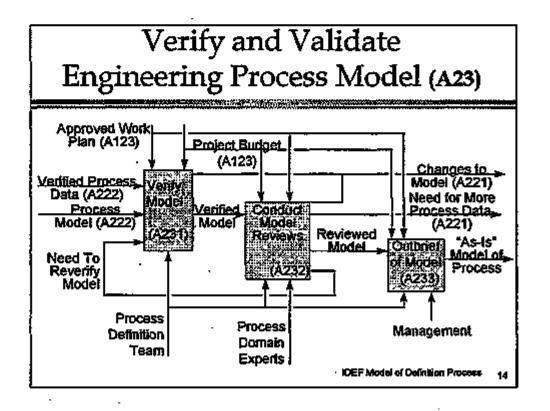


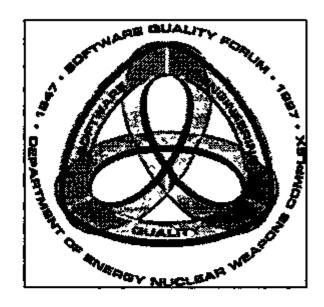












Session Y1: How the NWC Handles Software as Product

Chair Dave Vinson

Pantex Plant Software Quality Assurance Subcommittee Work Item #16

How the NWC Handles Software as Product Software Quality Assnrance Subcommittee Work Item #16

Presenters: Member(s) of SQAS WI#16: Management and Control of Product Software **Summary:** This tutorial provides a hands-on view of how the Nuclear Weapons Complex project\$ should be handling (or planning to handle) software as a product in response to Engineering Procedure 401099. The SQAS has published the document SQAS96-002, "Guidelines for NWC Processes for Handling Software Product," that will be the basis for the tutorial. The primary scope of the tutorial is on software products that result from weapons and weapons-related projects, although the information presented is applicable to many software projects. Processes that involve the exchange, review, or evaluation of software product between or among NWC sites, DOE, and external customers will be described. These processes include:

- 1. Identification: what are software product items, how are the product and items identified, how does software identification relate to system identification.
- Qualification: what is software qualification in accordance with EP401099, who is involved, how does a software Process Realization Team work, what is in a Qualification Plan and how does this Plan lead to a Qualification Evaluation Release.
- Acceptance: how does DOB accept software product, what is a Quality Assurance Inspection Procedure, how are product qualification and acceptance related, what are site and DOE roles, what is needed for customer use (interagency and external end-use).
- 4. Delivery: what is the mechanism for shipping and receiving software product, how is delivery accomplished between NWC sites, how is delivery accomplished between a site and external customer.

A Case Study of a recently completed project will be given to each participant for hands-on review of how the guidelines for handling software product have been applied. In particular, examples of project products used in the handling processes that will be reviewed include: Material List, Qualification Plan, Software Requirements, Test Plan, Maintenance Plan, Software Production Requirements, Traveler, Product-Disk Labels, Integrated Contractor Order, Certificate of Inspection, Shipper Label, Package Label, Complete Engineering Release, Qualification Evaluation Release, and Quality Assurance Inspection Procedure.

Site-specific issues and the tailoring of the handling guidelines for use in non-weapons applications will be discussed. Members from several sites who are on the SQAS WI#16 Working Group will be available to discuss the site-specific issues.

Hand-Out Material:

- 1. Totorial Slides
- 2. SQAS96-002, "Guidelines for NWC Processes for Handling Software Product," June 1996.
- 3. Case Study Notebook

Audience/Restrictions:

This tutorial is primarily intended for personnel who are or will be managing, developing or supporting software that will be delivered to or used by external customers. Tutorial participants must be a Department of Energy or Nuclear Weapons Complex employee. Although none of the material in this tutorial is classified, its content may be sensitive. A valid badge will be required for participants in this tutorial. If you have a question as to whether you can participate, contact a Forum representative.

Contact Information:

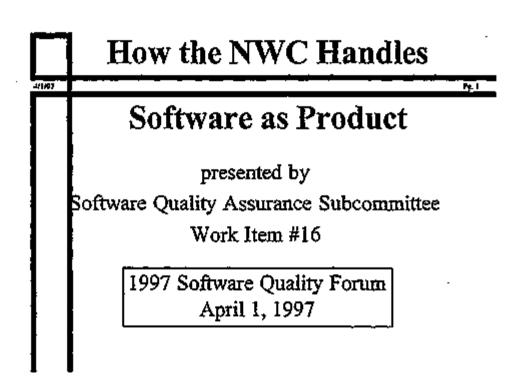
Dr. David E. Peercy Sandia National Laboratories P.O. Box 5800, MS0638 Albuquerque, NM 87185-0638 \$05-844-7965(voice), 505-844-3920(fax), depecte@sandia.gov

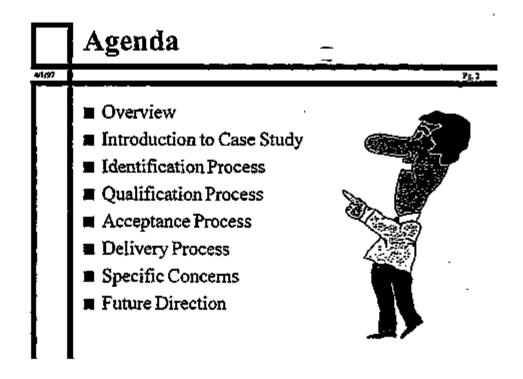
BIOGRAPHY

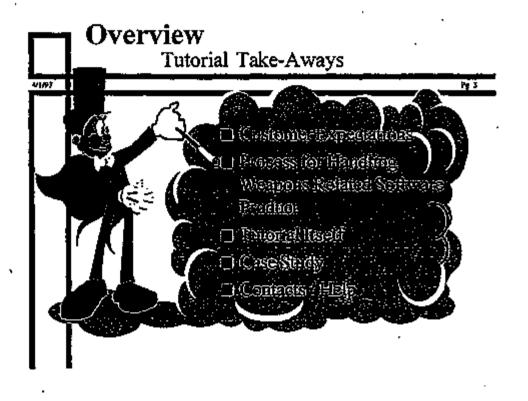
The Software Quality Assurance Subcommittee (SQAS) operates under the DOE Nuclear Weapons Complex (NWC) Quality Managers to identify and resolve Software Quality issues and problems common to all DOE sites and facilities. This tutorial is the result of an NWC SQAS work item to define how to manage and control software as product. The work item was established to satisfy a need to define a consistent process for handling product software. The Nuclear Weapons Complex-wide participants and presenters of this tutorial include:

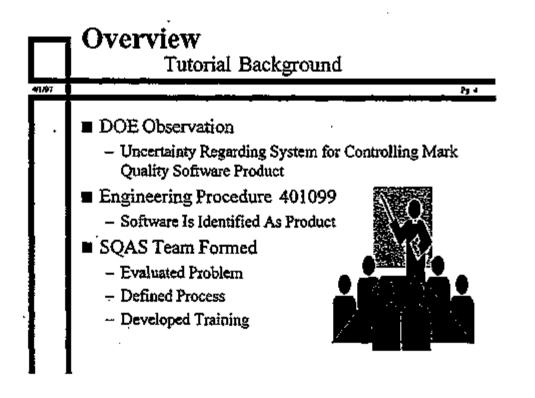
Chair David Vinson, Mason & Hanger, Pantex Plant Phil Huffman, Mason & Hanger, Pantex Plant Alvin Cowen, Mason & Hanger, Pantex Plant Catherine Kuhn, AlliedSignal Aerospace, Federal Manufacturing & Technologies Donald Schilling, AlliedSignal Aerospace, Federal Manufacturing & Technologies Dave Peercy, Sandia National Laboratories Mike Blackledge, Sandia National Laboratories Orval Hart, Los Alamos National Laboratory John Cerutti, Los Alamos National Laboratory Bill Warren, Lawrence Livermore National Laboratory Charles Chow, Lawrence Livermore National Laboratory Ellis Sykes, Department of Energy, Kansas City Office Gary Echert, Department of Energy, Albuquerque Office Kathleen Canal, Department of Energy, Headquarters Ray Cullen, Westinghouse, Savannah River Site Faye Brown, Lockheed Martin Energy Systems, Oak Ridge, Y-12 Plant

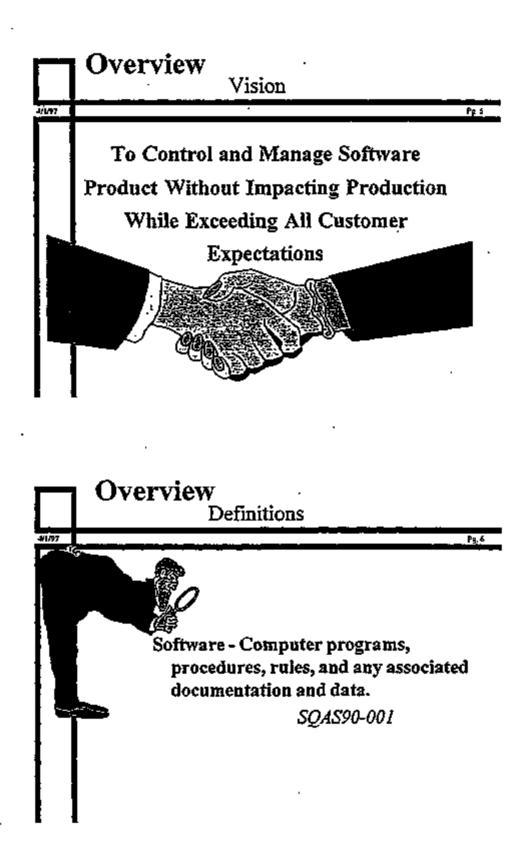
Dave Vinson, Chair WI#16 Mason & Hanger Pantex Plant Eldg 12-102 E.O. Box 30020 Amarillo, TX 79120-0020 Voice: 806-477-4739 Fax: 306-477-4350 E-mail: dvinson@pantex.com

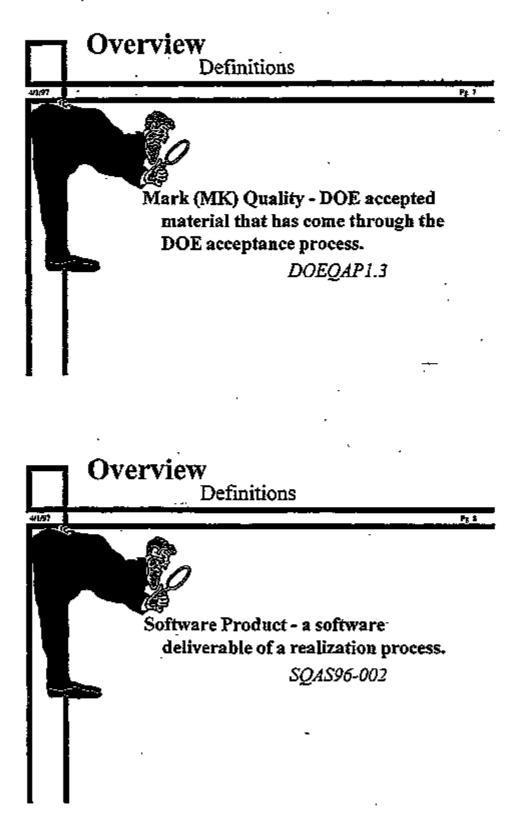






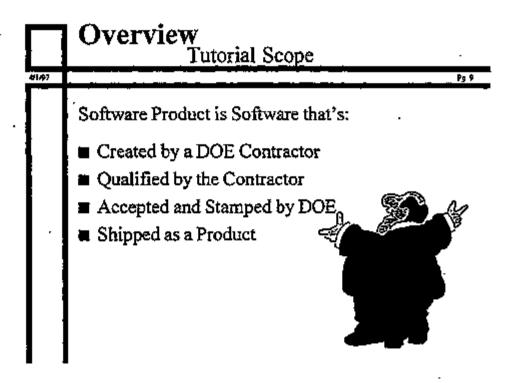


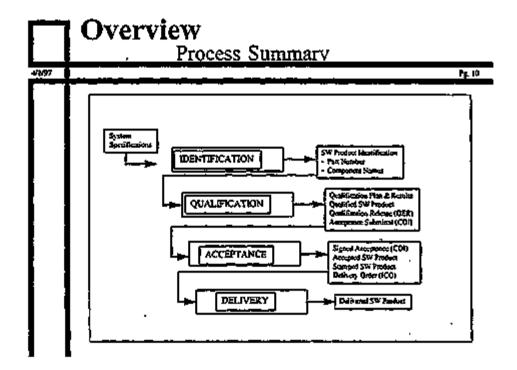


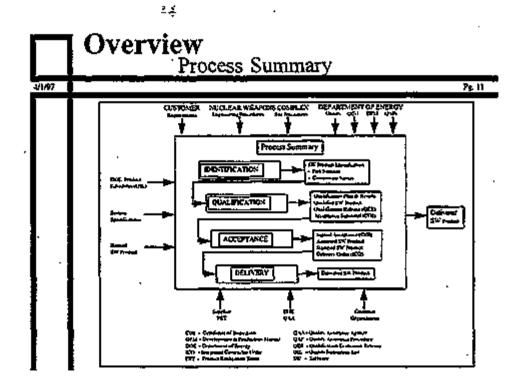


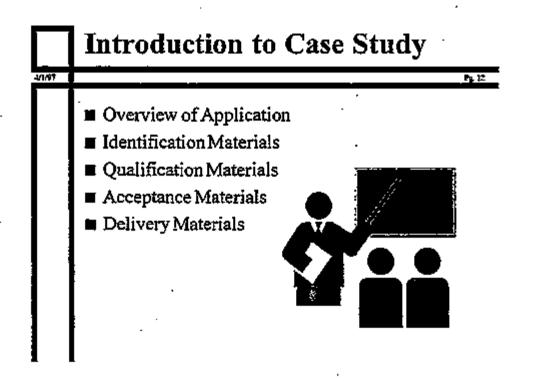
Page 4

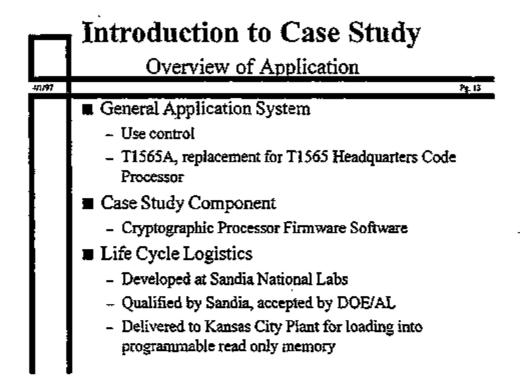
)

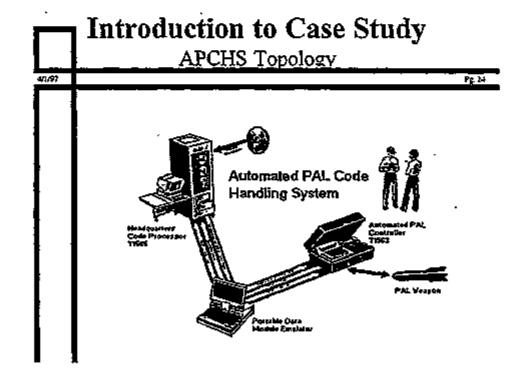


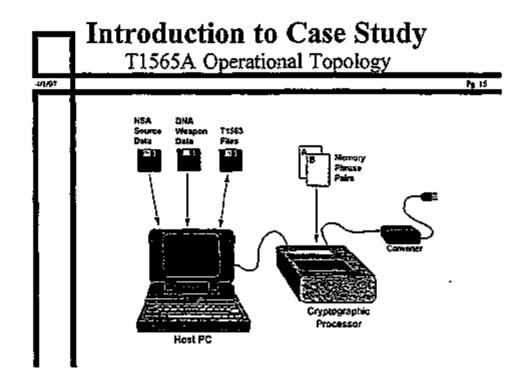


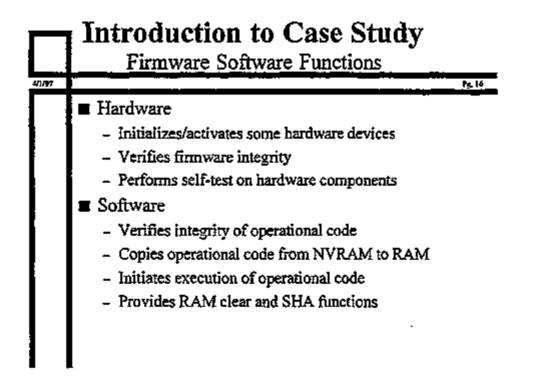


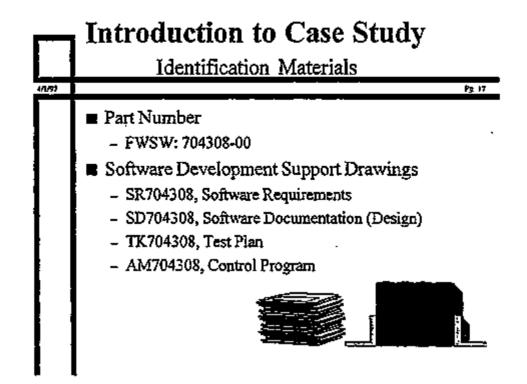


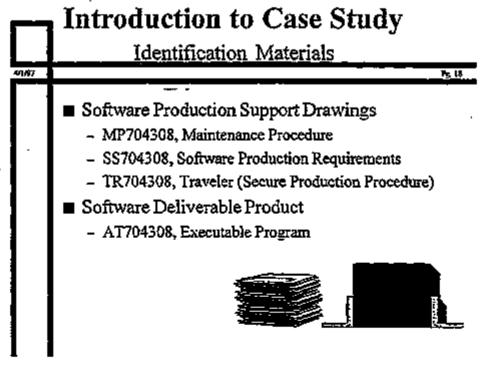






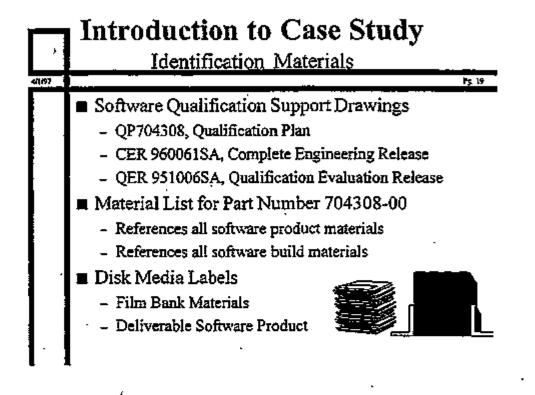


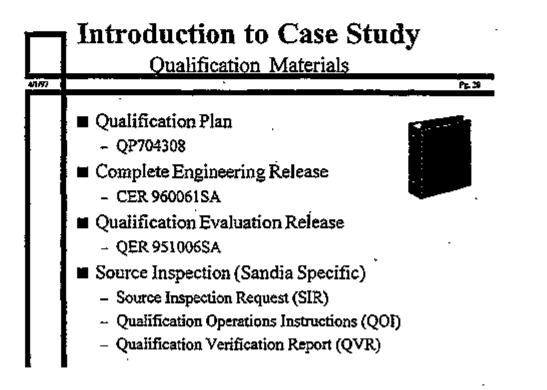


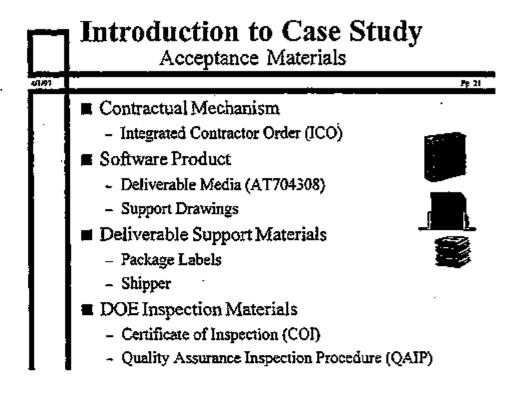


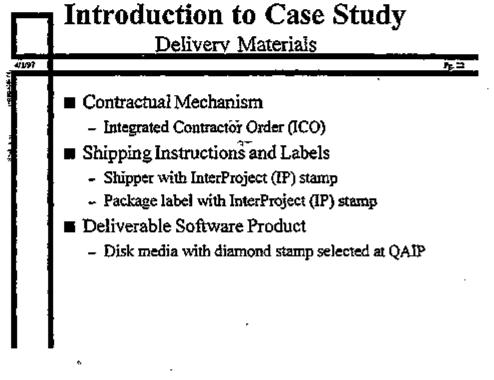
Ċ.

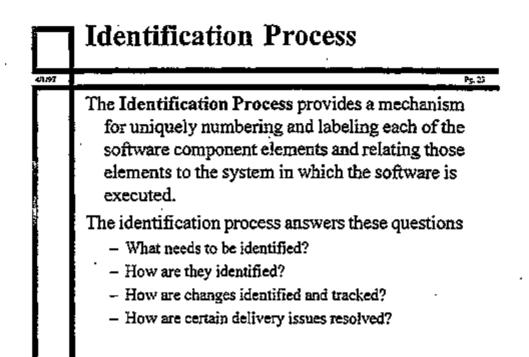
1997 Sofware Quality Forum 4/1/96

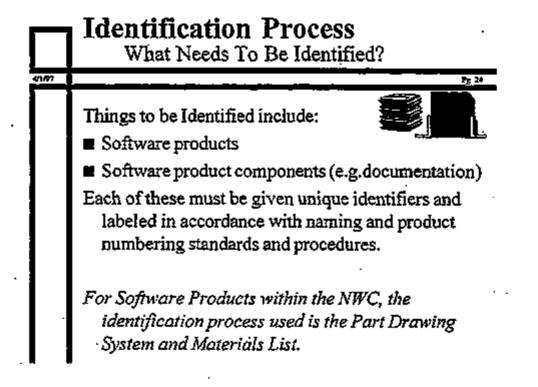


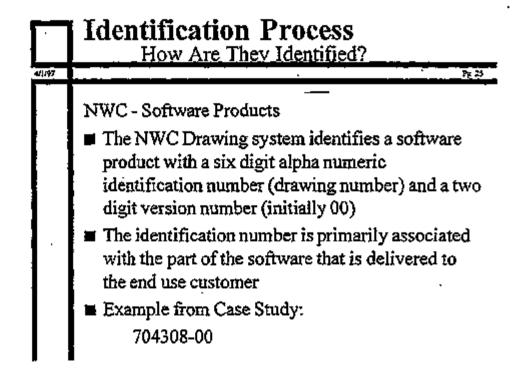




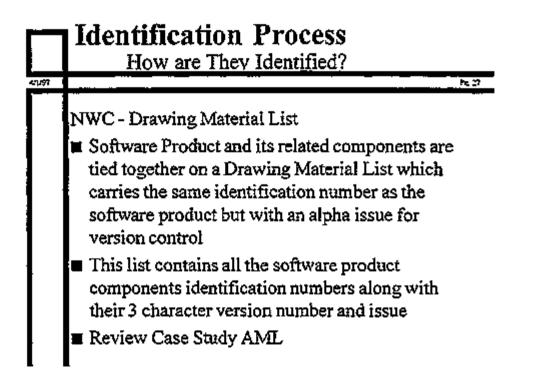


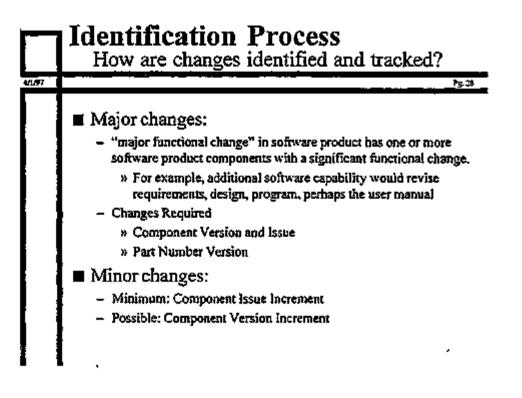


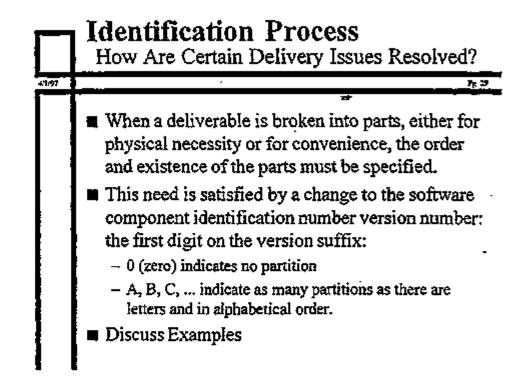


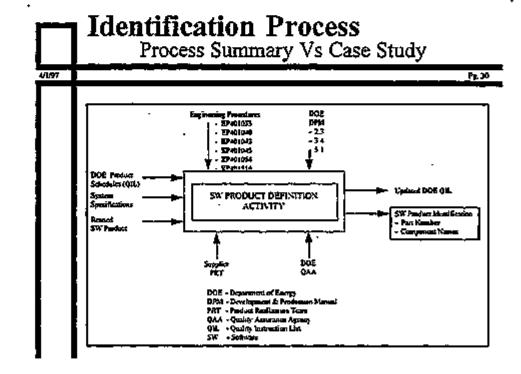


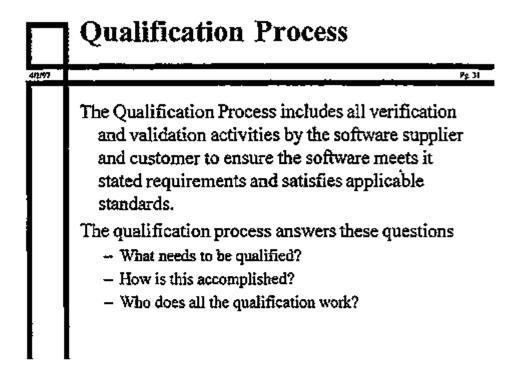
i	Identification Process How are They Identified?	
	Pg 36	
-	NWC - Software Product Components	
	All related software product components	
1	identification numbers are derived from the	
1	associated software product identification number	
	We do this by adding a 2 digit prefix indicating the software product component, a 3 character version number, and an alpha "Issue"	
	Example fröm Case Study	
	SR704308-000, Issue A	
	AT704308-000, Issue A (See Case Study)	
1		

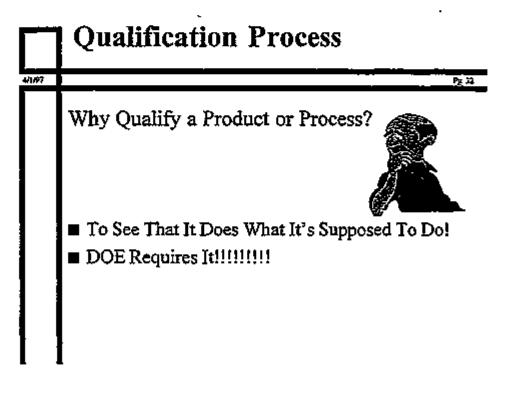


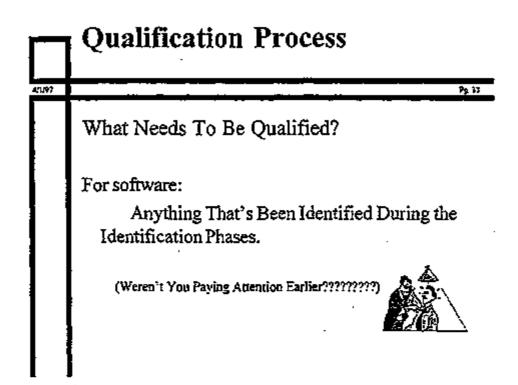




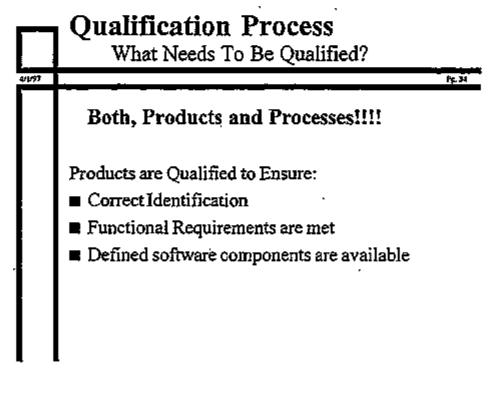




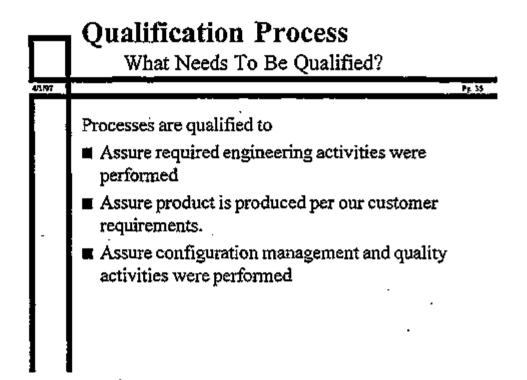


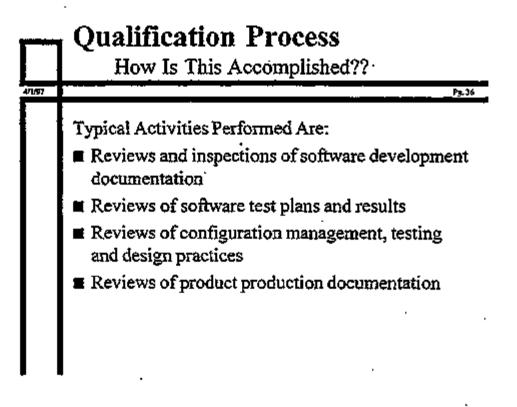


ī

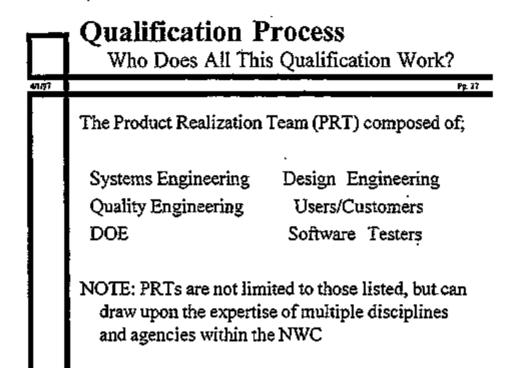


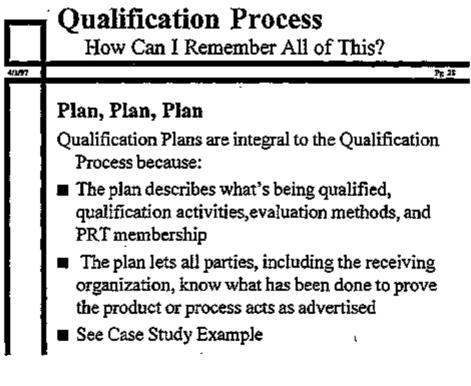
1997 Sofware Quality Forum 4/1/96

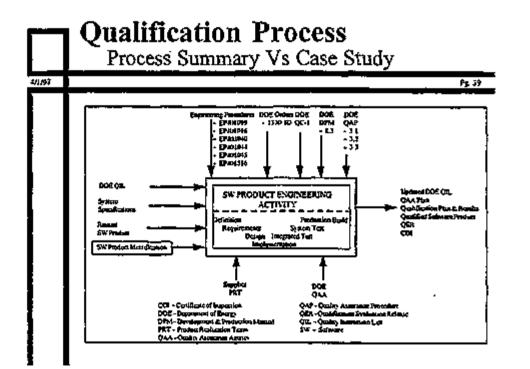


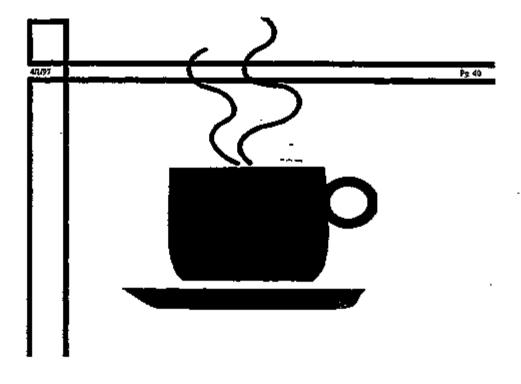


Page 18.





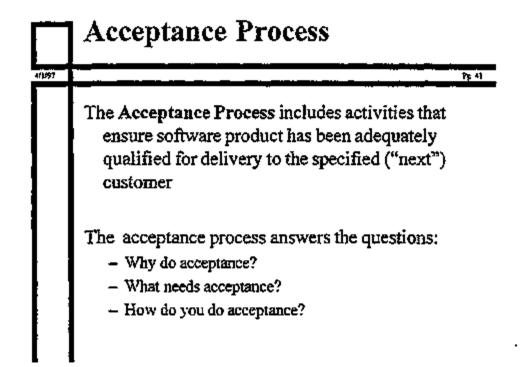


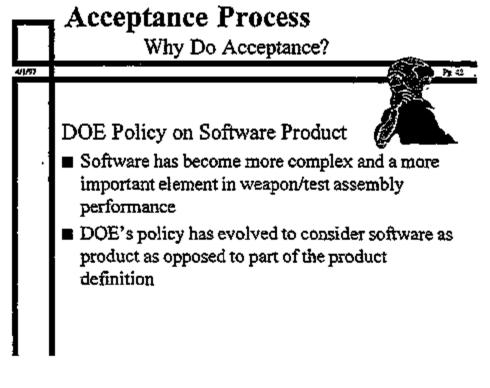


Page 20

.)

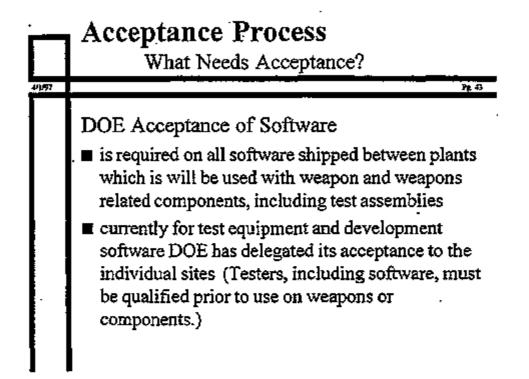
<u>_</u>)

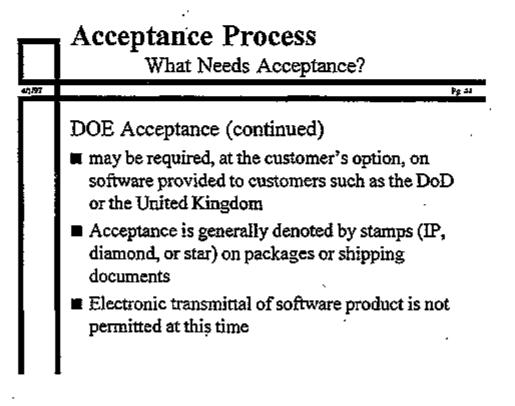




1

1997 Sofware Quality Forum 4/1/96





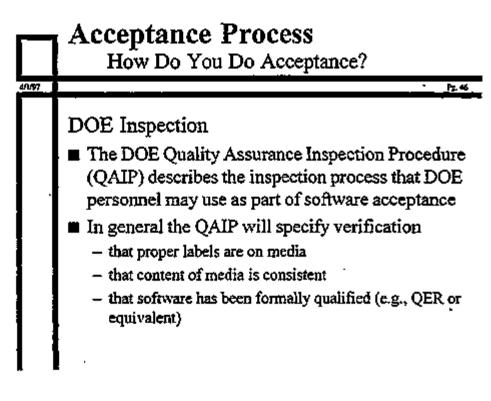
Acceptance Process

How Do You Do Acceptance?

Submission to DOE

41,197

The Certificate of Inspection (COI) is the form used by the contractor to submit software and other product to DOE, to identify the product definition requirements, and to certify that it meets those requirements.



1997 Sofware Quality Forum 4/1/96

P5 45

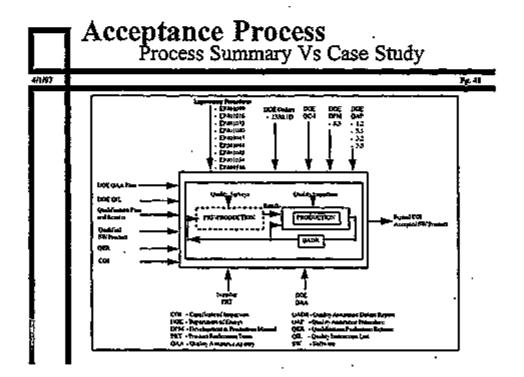
Acceptance Process

4/1/97

Summary

DOE Policy on Software Product

- Acceptance is essential for weapon software to provide an independent assessment that requirements have been met.
- The receiving agency requires an indication of DOE acceptance if software is intended for use in weapon product.



1997 Sofware Quality Forum 4/1/96

Page 24

74.57

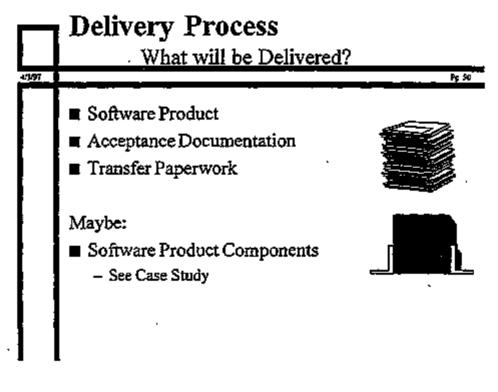
Delivery Process

4'3/97

The Delivery Process includes all supplier and customer logistic activities of shipping and receiving. The Delivery Process should be sensitive to the variations in delivery of developmental software product, prove-in software product, and production software product.

The delivery process answers the questions

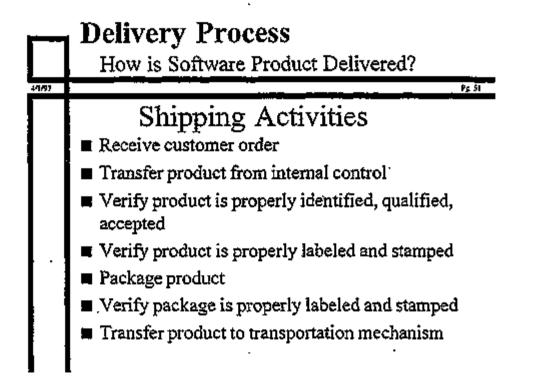
- What will be delivered?
- How is software product delivered?
- How are software product components delivered?

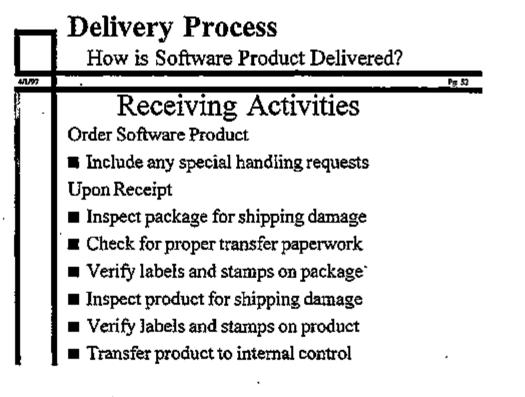


1997 Sofware Quality Forum 4/1/96

Page 25

Pg. 49





Delivery Process

Shipping Organization

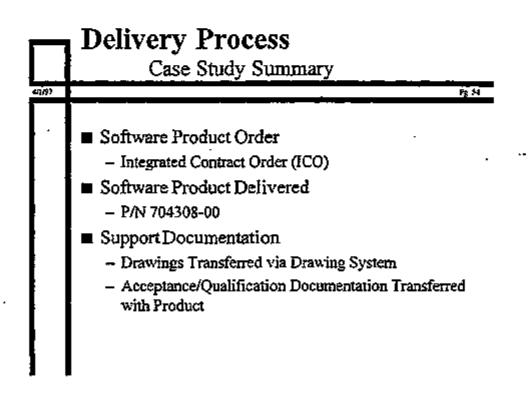
 Verify product acceptance documentation is complete

How are Software Product Components Delivered?

- Transfer any support documentation including
- drawings to the receiving organization

Receiving Organization

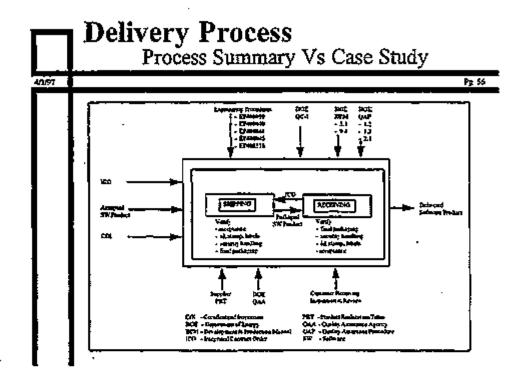
- Verify product acceptance documentation is complete
- Verify that support documentation is released and available for use

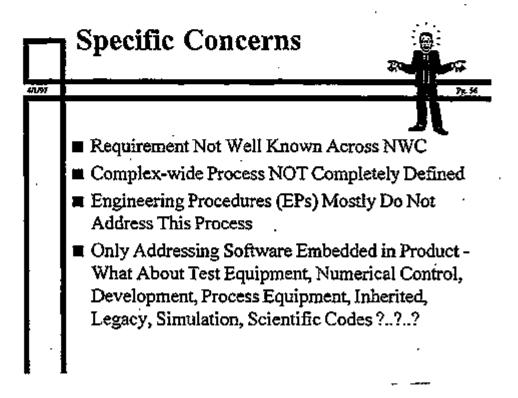


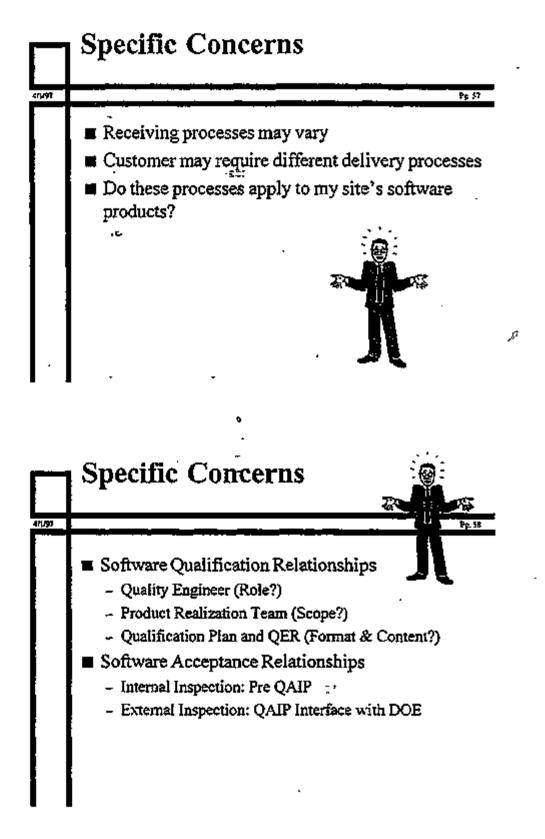
1997 Sofware Quality Forum 4/1/96

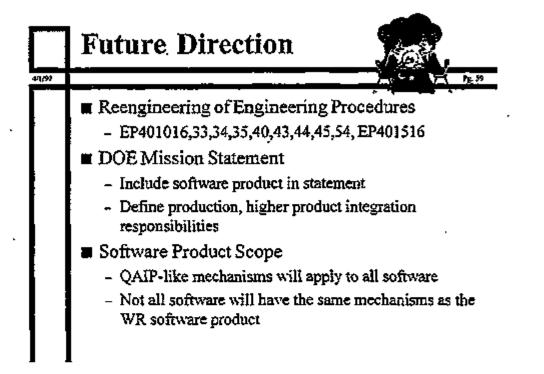
Pc 53

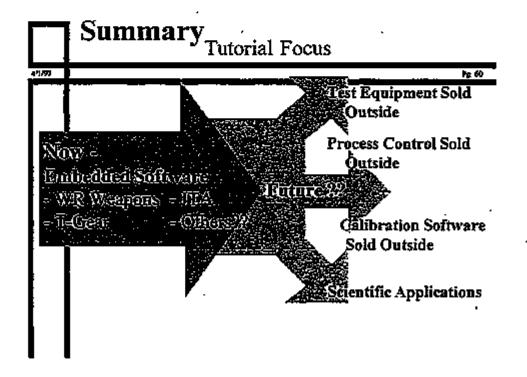
÷

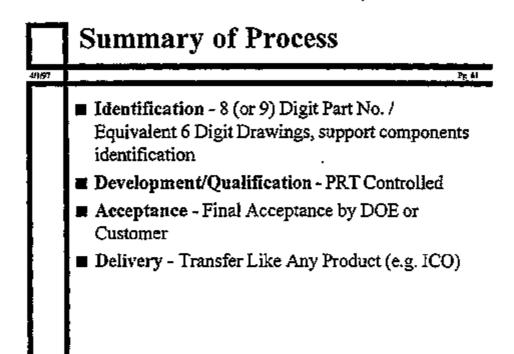


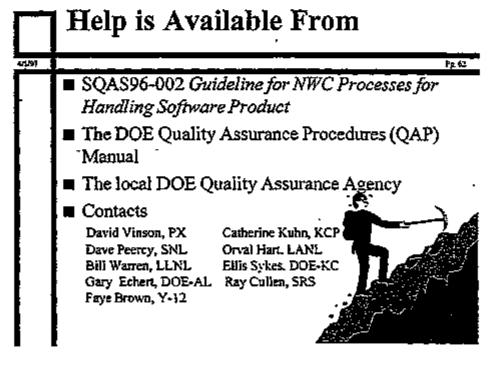














Session W2: Writing Testable Software Requirements

Dr. Dwayne Knirk Sandia National Laboratories

Writing Testable Software Requirements Dr. Dwayne L. Knirk Sandia National Laboratories

This tutorial identifies common problems in analyzing requirements in the problem and constructing a written specification of what the software is to do. It deals with two main problem areas: identifying and describing *problem requirements*; and analyzing and describing *behavior specifications*.

Software-intensive systems are expected to work in a particular environment to bring about desired effects in that environment. To accomplish these effects, the computing system must have a variety of interactions with that environment. Its capabilities and features are directed to establishing a variety of relationships between those interactions, including stimulus-response, constraint, and historical reference. To establish such relationships are the services provided by the computing system. The given environment and required effects in the problem are collectively documented as Problem Requirements. The computing system interactions and services are documented Behavior Specification. The relationship between these two sets of information is an explicit and verifiable behavior design task.

The Behavior Specification characterizes a computing system independently of its application context. Having a behavioral specification enables a true concurrence in development and testing processes. It provides a single reference point for all decisions of software architecture and implementation as well as for test case and testware architecture and implementation.

This tutorial focuses on determining what facts about a computing system are to be documented, how they should be expressed, and how they are related to facts about the application environment. It provides an overview of these basic specification techniques:

- the application of standard problem frames for classifying and organizing the various requirements,
- the application of stimulus/response and client/server viewpoints for structuring the description of computing system behavior,
- the expression of unique, testable action statements with the help of pre- and postconditions, state models, and datastore models,
- the description of behaviors of components and their architectural composition into the behaviors of assemblies, and
- the use of these descriptions in Software Requirements Specification documents.

Much of this material in this tutorial is being developed as part of the next revision of IEEE Std 1175, Part of that standard is a system behavior meta model. Various parts of the material are undergoing refinement by application in various Sandia projects.

BIOGRAPHY

Dr. Dwayne L. Knirk

Dr. Knirk is a member of the software quality engineering department at Sandia National Laboratories. He provides in-house consulting to line organization projects for software engineering processes, methods, standards, tools, and training. He participates in process assessments and improvement programs, and provides support for configuration management, software inspections, and process automation. Dr. Knirk's primary focus is on the two complementary areas of software specification and testing, in which he works to bring more formal methods into more practical applications. He works actively on IEEE software engineering standards groups. He is a member of the ASQC Software Division Methods Committee.

Dr. Knirk previously worked for Programming Environments, Inc., where he was the architect and principal developer of the automated software test design tool, T. That commercial product analyzed a formal software behavior description for testability, designed test cases for demonstrating that behavior, and generated actual test case data.

Knirk: Establishing a Three-Way Agreement

04/01/1997

Tutorial: Writing Testable Software Requirements

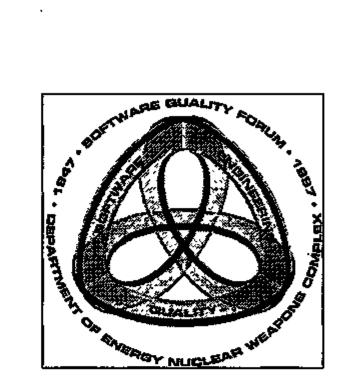
Software Quality Forum Albuquerque, NM 1 April 1997

Presented by Dr. Dwayne L. Knirk Quality Engineering Department Sandia National Laboratories, Albuquerque, NM

SAND97-XXXC

This work was supported by the United States Department of Energy under Contract DE-AC04-94AL85000. Sandia is a multiprogram laboratory operated by Sandia Corporation, a Lockheed Martin Company, for the United States Department of Energy

Sandia National Laboratories



Session X2: Using COTS Software in Development Projects

Q,

Lt Col Nancy Crowley Acting Chief, Space System Technologies Phillips Laboratory Kirtland AFB, NM

The Use of COTS in the Multimission Advanced Ground Intelligent Control (MAGIC) Program

Lt Col Nancy L. Crowley, Phillips Laboratory PL/VTS

The use of commercial software and standards has been touted as a potential for significant cost and time savings in developing military systems, specifically, satellite control systems. And while the savings do exist, commercial software and standards must be carefully evaluated prior to selection, carefully integrated, and used where appropriate to reap their benefits. For example, not all Commercial Off-The-Shelf (COTS) products are suitable because they encompass too may inseparable functions, have a very narrow customer base and/or have no possible replacement COTS products. A COTS-based system should consist of small components that do one contained task and integrate with other components through some sort of message passing, such as files, DDE, OLE, DLL or other appropriate middleware protocols such as provided in the CORBA environment. A component should be able to be replaced with no, or minimal, impact on other components in the system. Commercial protocols can be unstable and change rapidly over time, forcing decisions on when to upgrade the components to new versions, and evaluating the impact of doing so. Also, COTS components have bugs, and are usually not tested to the stringent standards seen for some military systems. The features in COTS components are often not exactly what is needed, necessitating decisions on whether they are good enough, or if some custom code should be developed and integrated.

The tutorial will discuss the experiences of the Space System Technologies Division of the USAF Phillips Laboratory (PL/VTS) in developing a COTS-based satellite control system. The system's primary use is a testbed for new technologies that are intended for future integration into the operational satellite control system. As such, the control system architecture must be extremely open and flexible so we can integrate new components and functions easily and also provide our system to contractors for their component work. The system is based on commercial hardware, is based on Windows NT, and makes the maximum use of COTS components and industry standards.

BIOGRAPHY

Nancy L. Crowley, Lt Col Acting Chief, Space System Technologies

÷

ł

Ł

1

ł.

L

I.

Lt Col Nancy Crowley is the Acting Chief of the Space System Technologies Division (PL/VTS), Kirtland AFB, New Mexico. The focus of Space System Technologies Division is on the innovative application of software technologies to improve performance and reduce operations and maintenance costs for satellite control systems, including telemetry, tracking and commanding (TT&C), mission data dissemination, data processing, and satellite autonomy. Lt Col Crowley is also the program manager for the Multimission Advanced Ground Intelligent Control (MAGIC) program. MAGIC is developing the architecture for the next generation satellite control system that provides a low cost, flexible software architecture that allows plug and play of COTS products in a vendor independent manner.

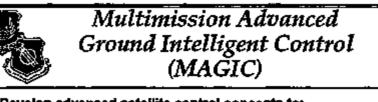
Lt Col Crowley was born May 13, 1955 in the Bronx, New York. She graduated from Theills High School in Theills NY, in 1973. She received a Bachelor of Science in Electrical Engineering from the University of New Hampshire in 1977 where she was a ROTC distinguished graduate. She later received the Master of Science in Digital Engineering and the Doctor of Philosophy (major of software engineering, minor of artificial intelligence) from the Air Force Institute of Technology in 1982 and 1994 respectively. Her research was in object-oriented methods for software requirements analysis.

Lt Col Crowley entered the Air Force in 1972 and was a flight test engineer for Tactical Air Command. There she conducted operational test and evaluation and flew in fighter aircraft in support of projects. After her masters degree, she was assigned to the Flight Dynamics Laboratory, where she was the software engineer for the digital flight control system of the X-29 Advanced Technology Demonstrator and the Ada focal point for the laboratory. There and in subsequent assignments she was a technical consultant to the Swedish government on the development of the digital flight control system for the JAS-39. Her next assignment was at the Systems Acquisition School, Brooks AFB Texas where she was a course developer and instructor of software acquisition courses. There she was also a system administrator for a UNIX and PC-based networked system that serviced the students and staff at the school. After completing her Ph.D., she came to her current assignment in Oct 94.

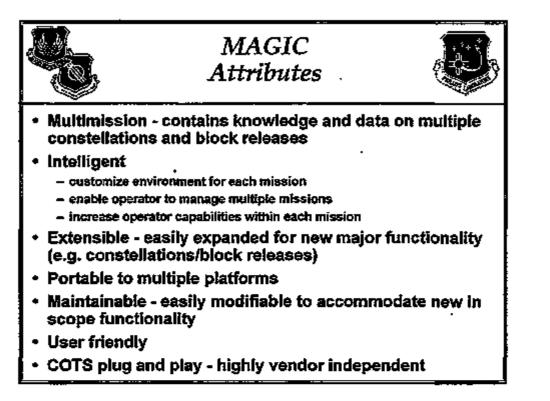
Outside her Air Force duties, Lt Col Crowley teaches software engineering, software management, and computer science courses at local Universities. Her and her husband own a computer consulting business. Both her and her husband enjoy riding horses.

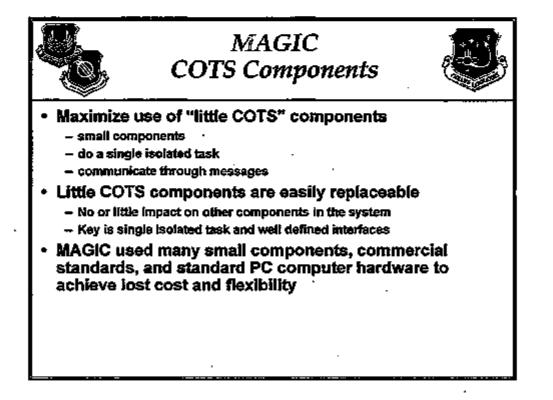
Phillips Laboratory PL/VTS 3550 Aberdeen Ave SE Kirtland AFB NM 87117-5776 Voice: 505-846-0461, ext 313x: 505-846-6053 ...mail: crowleyn@pik.af.mil

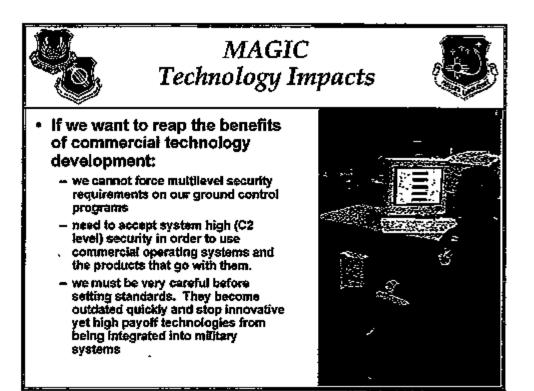


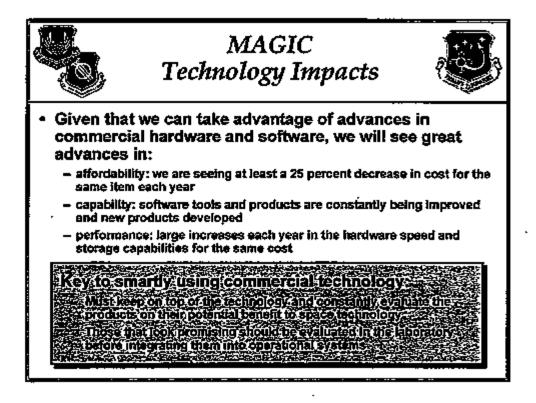


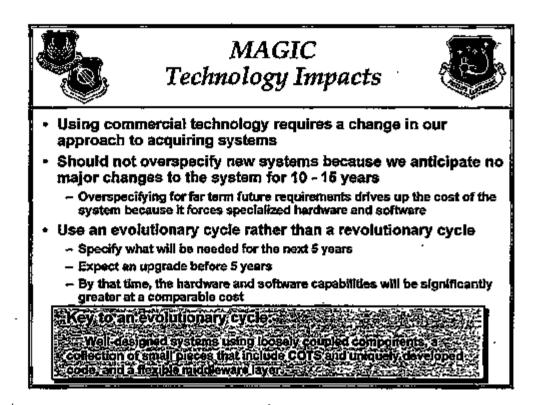
- Develop advanced satellite control concepts to:
 - Improve operator effectiveness
 - Support new ops concept (front room/back room)
 - Enhance operational capability
 - Reduce USAF Satellite Control Network (AFSCN) costs
- Focus on:
 - Telemetry analysis
 - Decision support
 - Operator training
- Integrate technology into USAF core TT&C system
- Technology: Use COTS integration, message passing between components, open distributed systems, object-oriented development, relational and object-oriented databases, and automated reasoning techniques to develop the next generation ground stations.

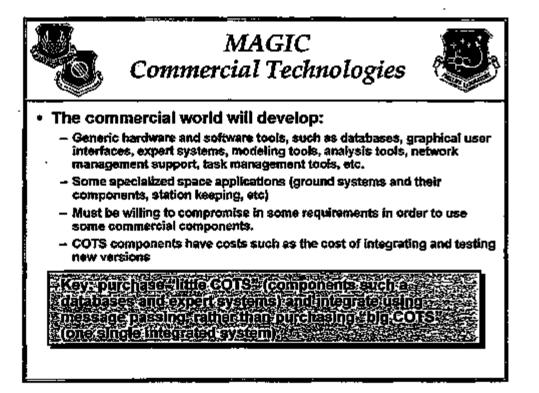


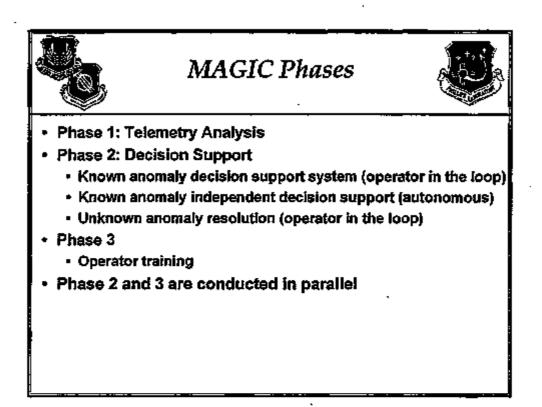


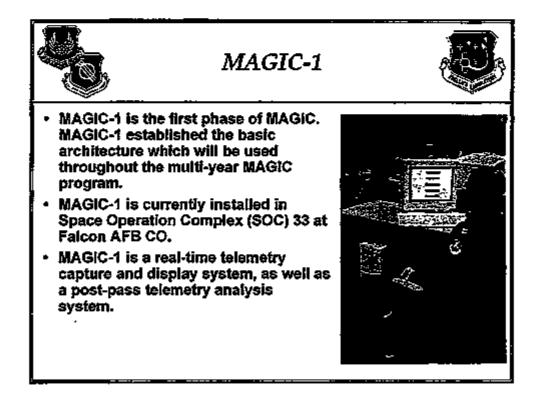


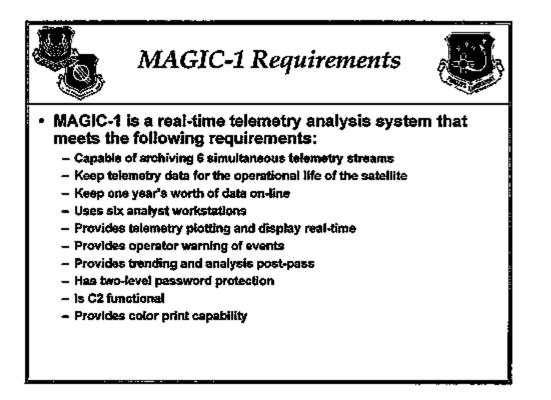


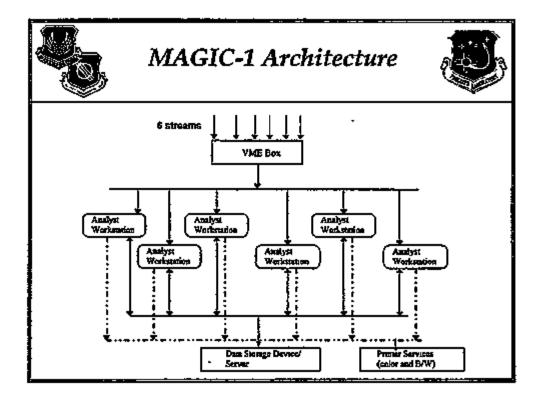








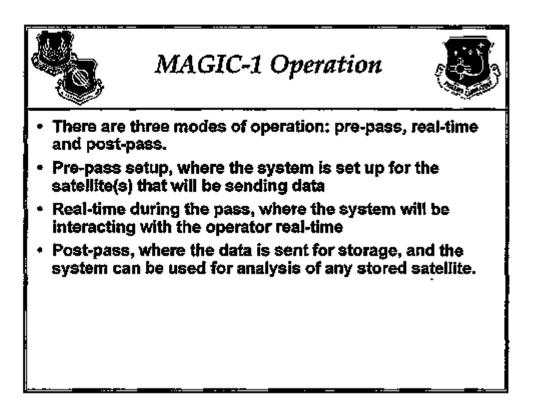


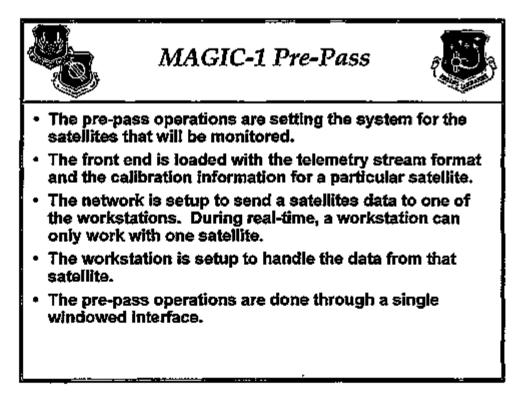


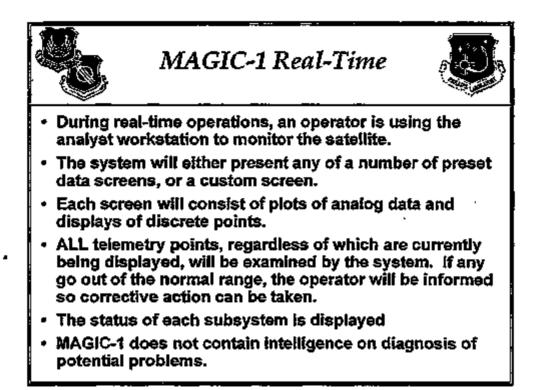
I

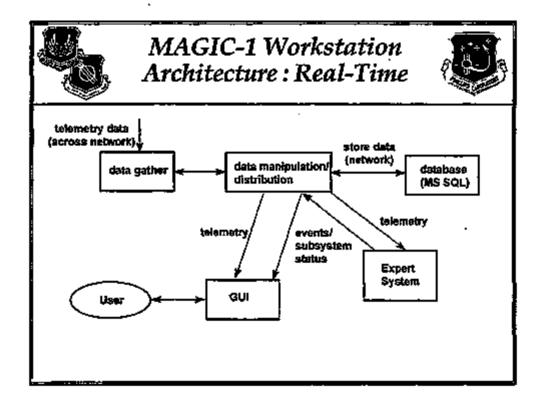
L

L



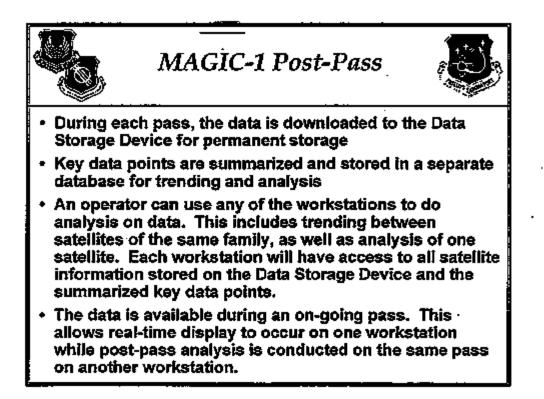


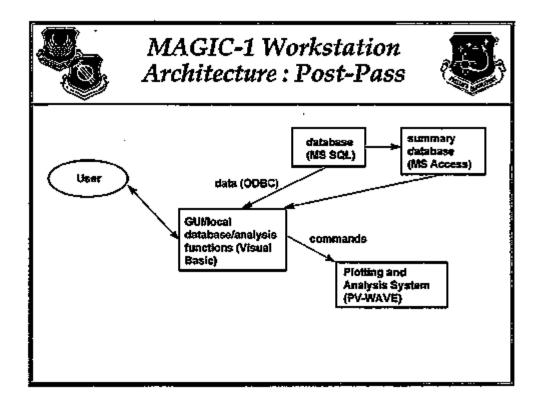


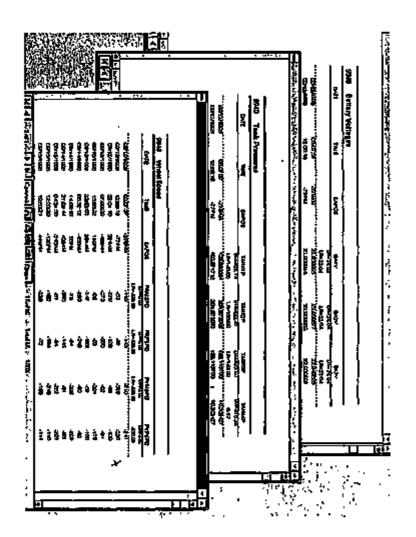


		}						SCN een			9 		
6275	13	768	7 0 7	50	77	2051+17	2	•77 EN	162	COMPOSITR	ATTON	VSNO	12 AD 10
ACE	в	OH	FR+MEZ.		CFR.	25.051	13	UPDATE	165BC			DHIA	10120100
ACCORP		DIS	PX-NRL		CO1	1.5 0157		5180.00	ENG.	REPORT A	OH:	0012	01001010
TIME	3		PY-MIL		00	YSER59	1/10	PATCH	DEX				
ACECIE	в	сы,	FY+NEL					EADED	22	IBPOID RS		1112A	02000010
MET N	ъ	010				ZPLOG	2294	BIASUP	DØ.	HOXOD RE			
RAM	A	08	23	_	CH	111516		RACCUP	200	NEXCHO RS	1.1		
XCH	2	0.4	ESOFY		20							DH3A	00022011
ИСТА/Ю	8		ESSEL		в	TERLOC		SUBACO	2/8	NEU-CE	A	DITSB	10001010
						728060		EARTERLY		NUNCH	A.		
100.07	1	ON .	R+55		-	CHIEF!		LONGCL	100	H2XCB	x	2WLA	00201101
P21		a.	R-35					TATLE OF	DIS			DHER	20200001
2014.H 1 1514.MAP+	2345	500 300	₽+55 ₽+66		B B	e/weire NSOFF7		nscord 11xpnf	DIS DIS	GDAPORT (GDACND RS		DHIC	12100211
												DWSA	00000000
10110	0	900	550200		DI	S PIER	K 659		A 201			DHSE	00000000
						1111	A		A ER			DHEA	00000000
2202		999	530701		P	TTHE	A 604		Y ER			0110	01000100
90.09		PR.	GIRO		OF	P NSTRR	LIS	NECRE	DIS	GDARLC .	1.37		

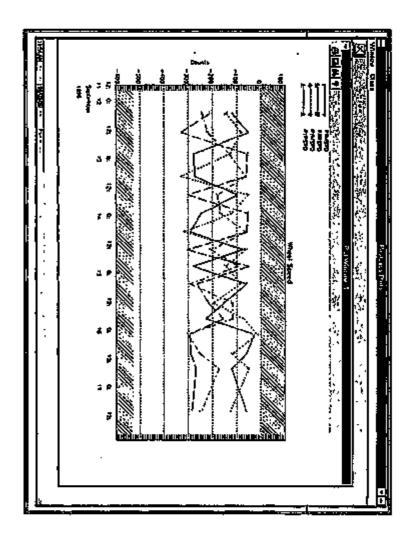
.	Visial Internet in SAM of Transmiss At Mysica
	Participation of the second state
4- 100 3 - 2720	Press and the second and second a
87-86 A.C.483	「「「「「」」」「「「」」」」「「」」」」」「「」」」」「「」」」」」「「」」」」
The second state	
10.00	
Real Street	
C SAHER SAL	
PENT -	
1.00	Line and the provide and the second states of the second states and the second states an
	See The second structure and shares
<u> </u>	
11.5	
H.	
Covel-Dates	
	the state of the s
1. 10 A	
	the second se







(!

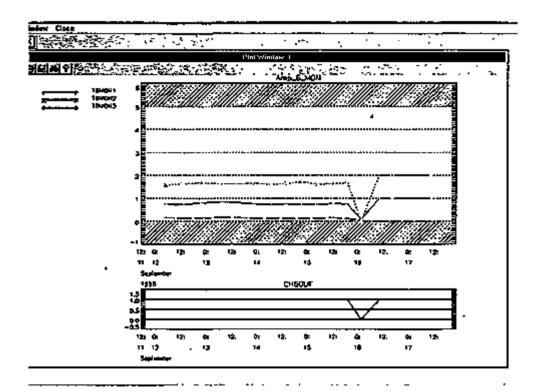


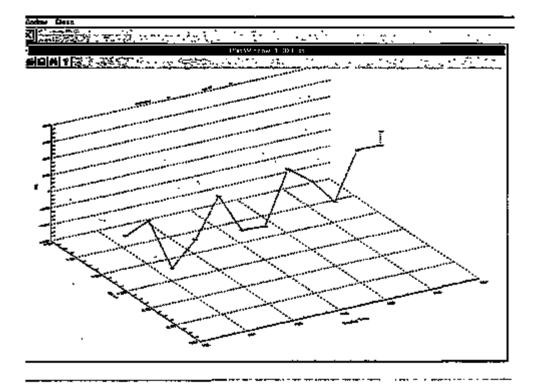
Page 11

1 1

1

L



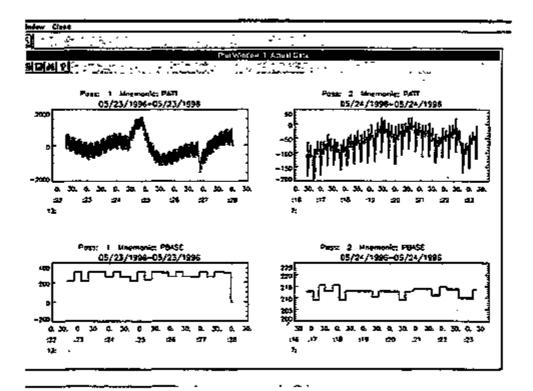


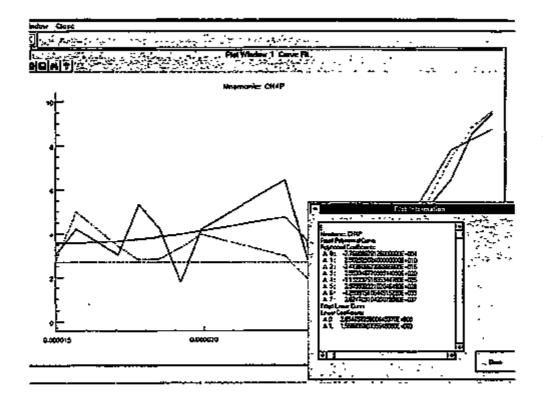
20

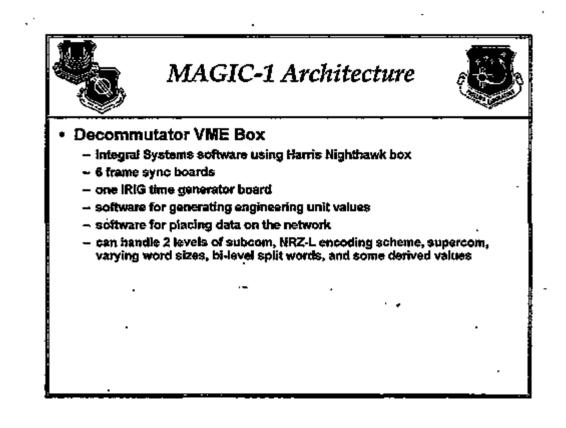
:

ı.

ົ)



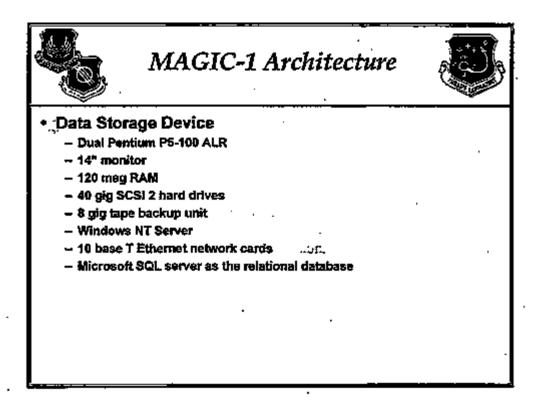


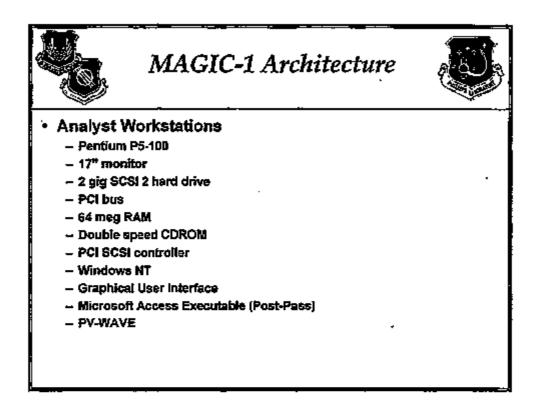


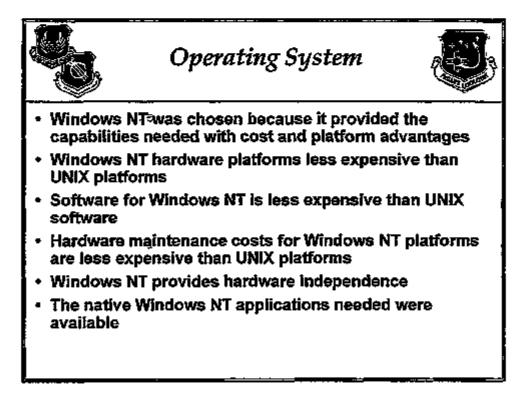
÷.,

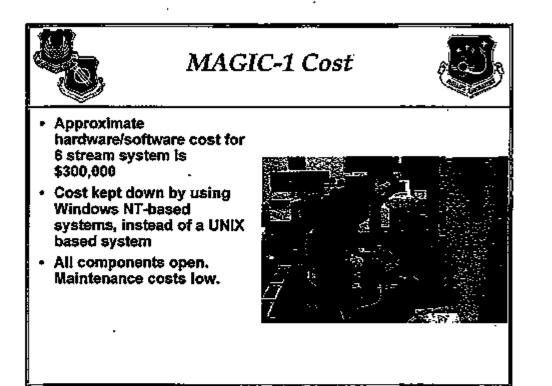
Ξ.

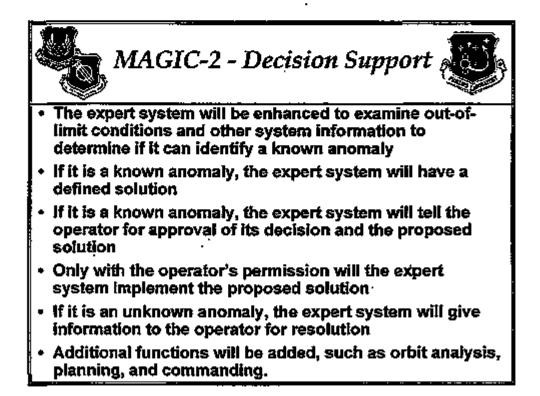
÷

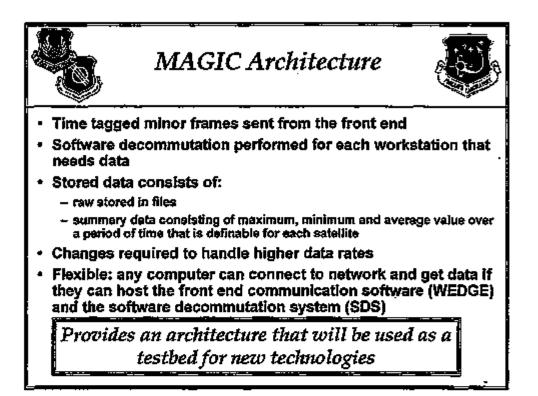


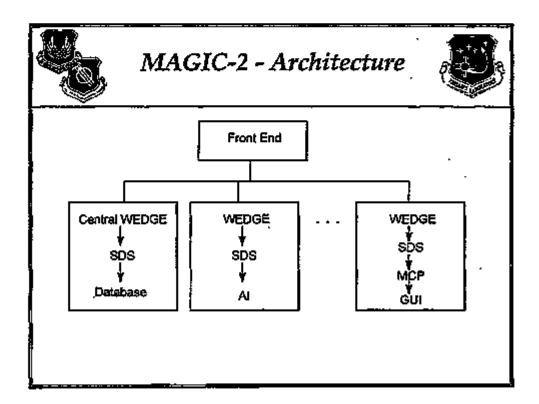


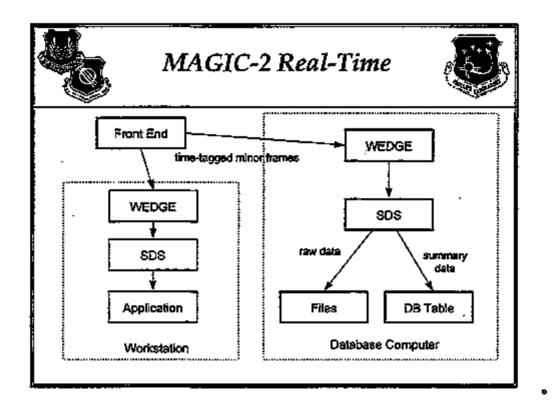


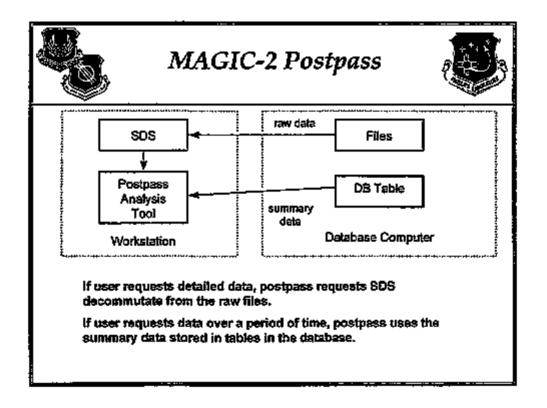






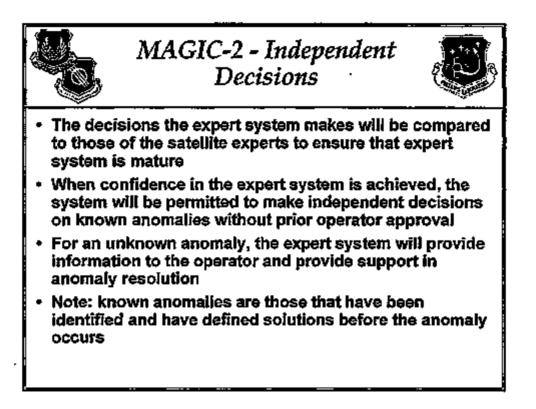


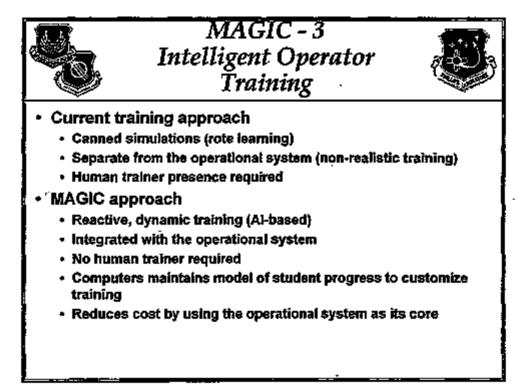


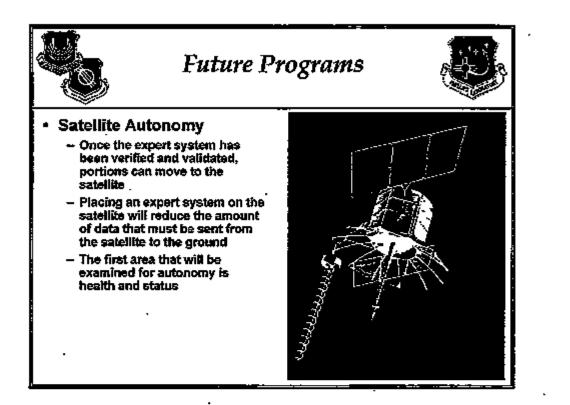


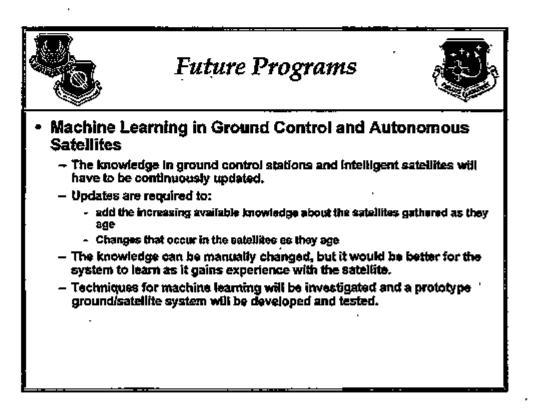
ı.

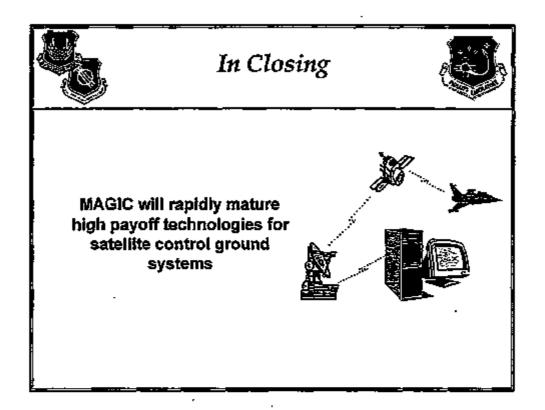
ı.













Session Y2: Software Inspection Process Overview

Larry Lane & Randy Dabbs Sandia National Laboratories

Overview of the Software Inspection Process

G. Lawrence Lane and Randy Dabbs Sandia National Laboratories

The Software Inspection is a formal in-process review method that provides immediate improvement in software product quality and produces metrics that indicate opportunities for process improvement. When adopted as a part of a defined, repeatable software development methodology, Software Inspections provide a mechanism for process control. The Software Inspection Process is not limited to formal reviews of code but applies to all software products. Software Inspections have consistently been shown to be very cost effective and is one of the most efficient ways to remove defects in all software products.

This tutorial introduces attendees to the Inspection Process and teaches them how to organize and participate in a software inspection. The tutorial advocates the benefits of inspections and encourages attendees to socialize the inspection process in their organizations.

The processes which are introduced in this tutorial agree with the methods recommended in the Sandia Preferred Processes for Software Development.

BIOGRAPHIES

G. Lawrence Lane

Larry Lane is a Senior Member of the Technical Staff at Sandia National Laboratories. He earned a Master of Arts Degree in mathematics from the University of Kansas. Larry joined Sandia Corporation in 1959 as an assembly language programmer in the field data reduction department. He has also worked as a operating systems programmer and was responsible for the selection and installation of Sandia's first general purpose time sharing computer. Larry also worked as a computer consultant for large scientific computers, as the second computer ombudsman, and was responsible for the development of an electronic tracking system for electrical testing of radiationhardened microcircuits.

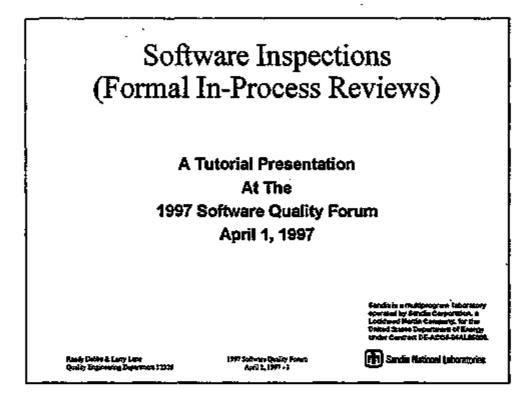
Larry moved to his current position in the Quality Engineering Department in 1991, where he is an instructor for the Software Quality Engineering course and the Software Inspection Class. As a software quality engineer, Larry has led numerous qualification efforts for new and upgraded software projects, particularly in the areas of use control and weapon security. He has helped develop and teach a customized version of the software inspection course to meet specific Sandia organizational needs.

Randy Dabbs

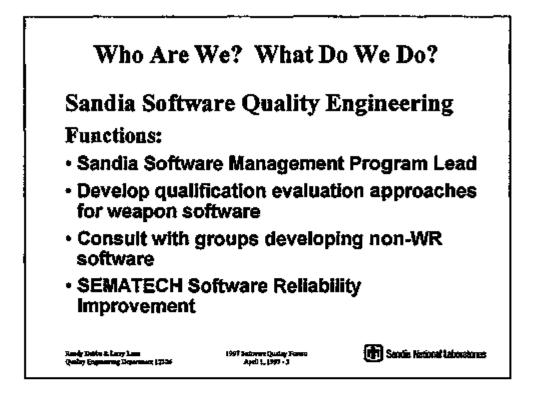
Randy Dabbs is a Senior Member of Technical Staff at Sandia National Laboratories. He has earned a Master of Science in Electrical Engineering from the University of New Mexico. He has beld positions at the Sandia Particle Beam Fusion Accelerator in the areas of data acquisition and signal processing; the Kwajalein Missile Range in the areas of range computer systems engineering, range operations, tracking software modeling and development, reentry mission project engineering, digital radar signal processing, radar controller real time software, and software configuration management; and the Sandia Kauai Test Facility in the areas of range computer support and operations, range safety software development, countdown software development, CASE tool selection and modeling of range operational software.

In his current position with the Sandia Quality Engineering Department, he has participated in instructing the Software Quality Engineering course and the Software Inspections course. In his role as software quality assurance engineer, he has participated in numerous software inspections for both internal and external customers. In addition, he has helped develop and teach a customized version of the software inspection course to meet the specific needs of Sandia organizations.

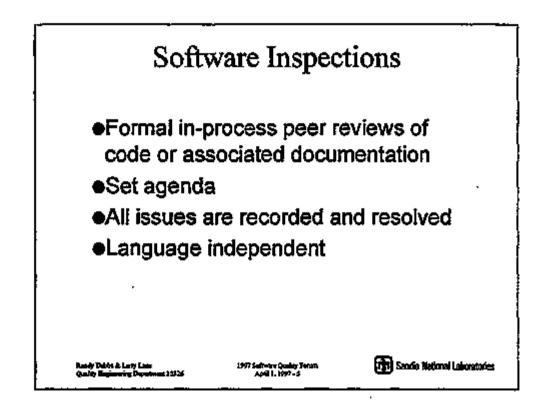
G. Lawrence "Larry" Lane Sandia National Laboratories PO Box 5800, MS0638 Albuquerque, NM, 87185-0638 Voice: (505)845-9122 email: gllane@sandia.gov Randy Dabbs Sandia National Laboratories P.O. Box 5800 MS-0638 Albuquerque, NM 87185-0638 Voice: 505-845-9232 email: rddabbs@aandia.gov

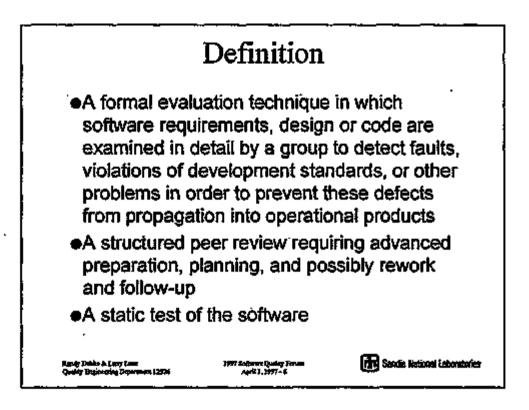


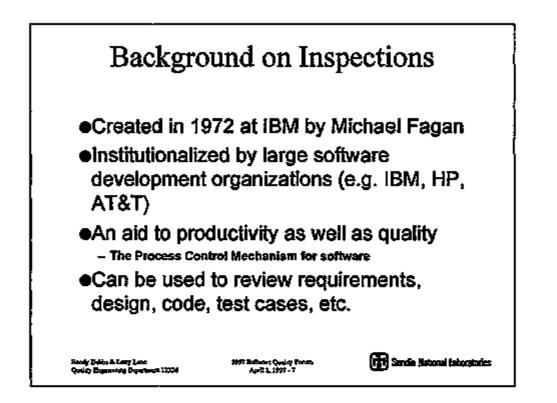


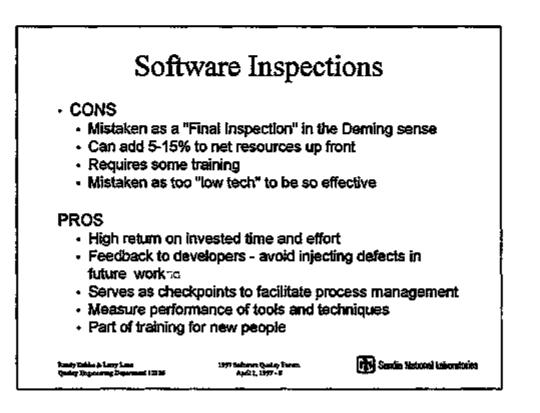


Т	utorial Goal	S
Introduce the	Inspection Proces	is is
- Learn how to o	rganize and participate i	in inspections
- Understand the	a major elements of soft	ware inspections
> Participant I	Roles	
> inspection i	Process Steps	
» Guidelines (ior Effective Use	
 Experience the 	Inspection process thre	sugh the workshop
 Socialize the l 	inspection Proces	5
- Recommend at	itendence at a formal ins	pection course
- Recommend in	nspections on your soft	ware products
Advocate the	benefits of Inspec	tions
- Cost savings	•	
- Shorten deliver	ry schedule	
- Reduction in de	efects	
Annaly Diskin & Lang Lang	till fallense finder fan en	🕅 Santin Katoral Laboratorias

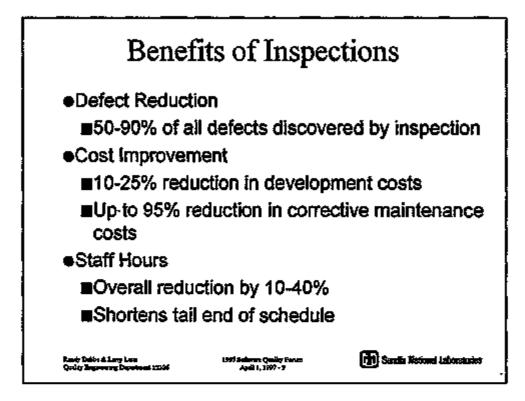


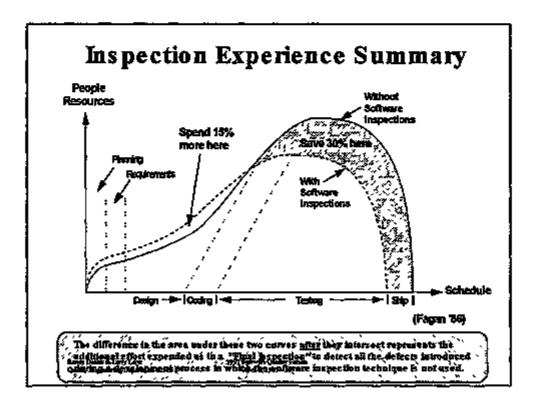


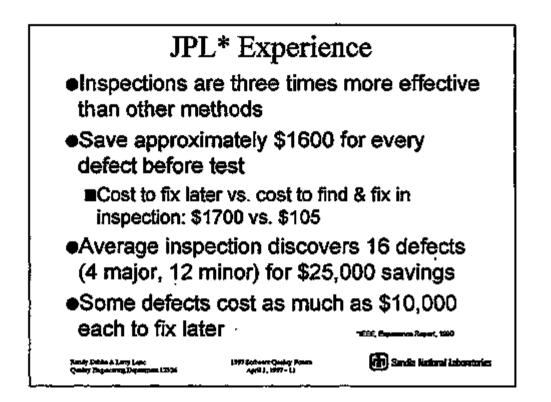


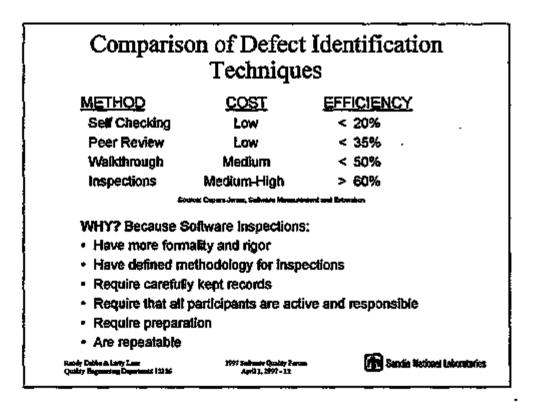


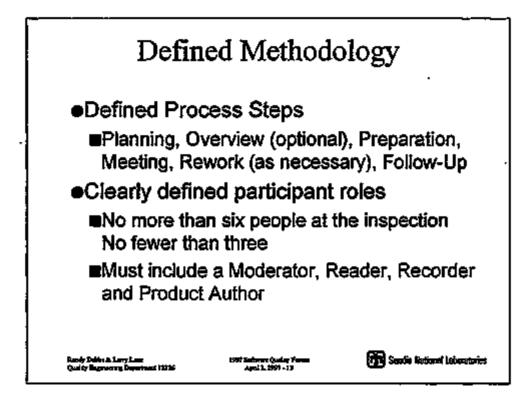
ι

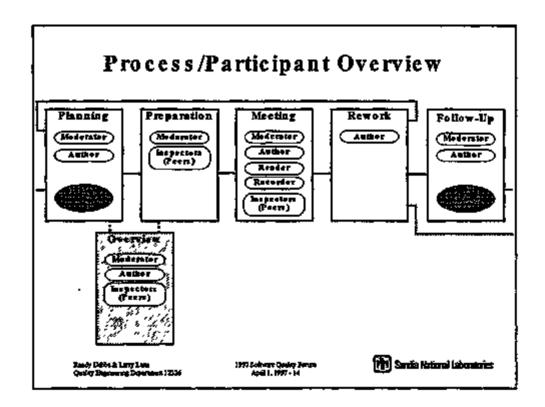


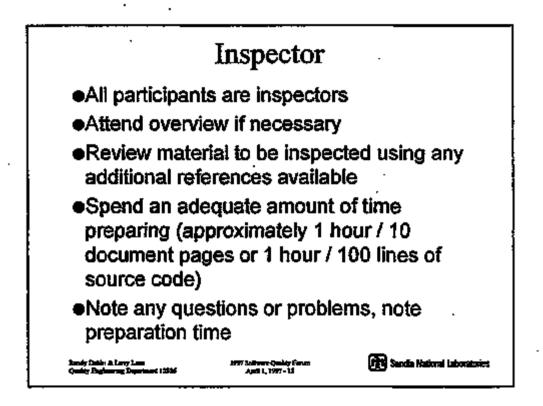


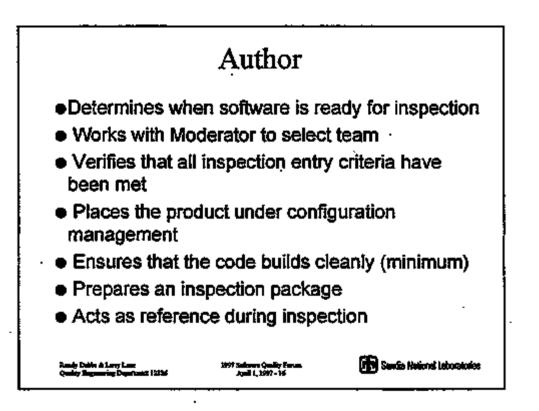


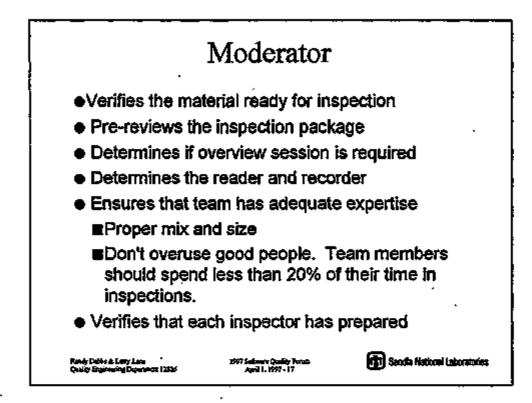




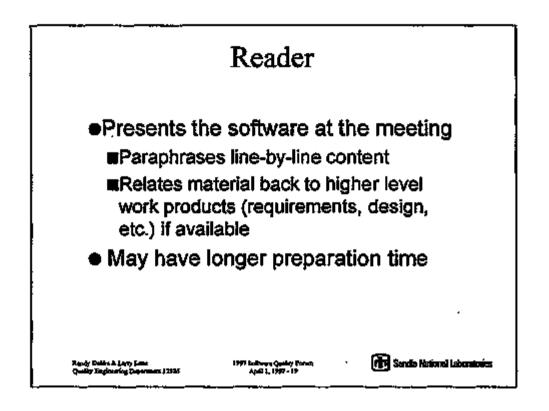


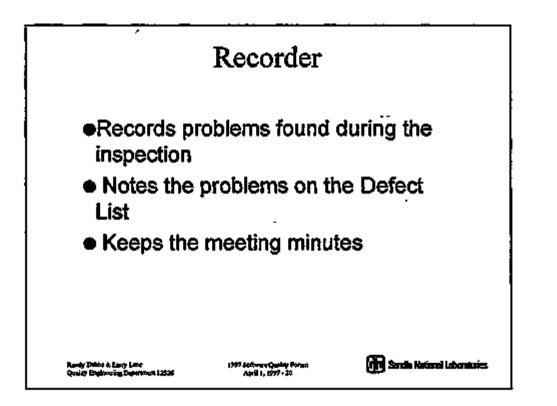




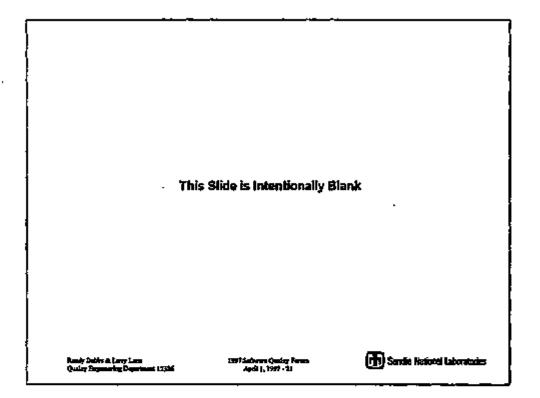


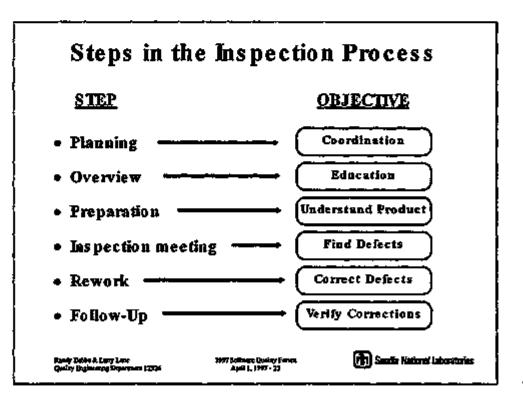




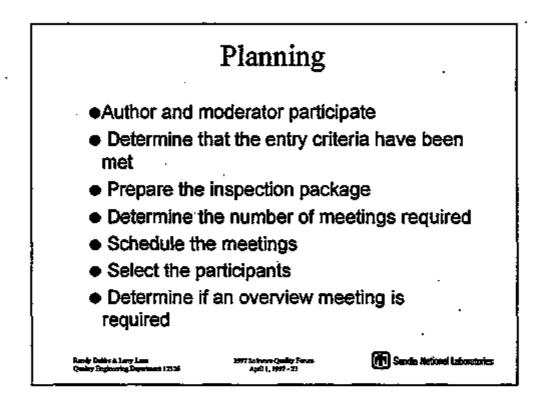


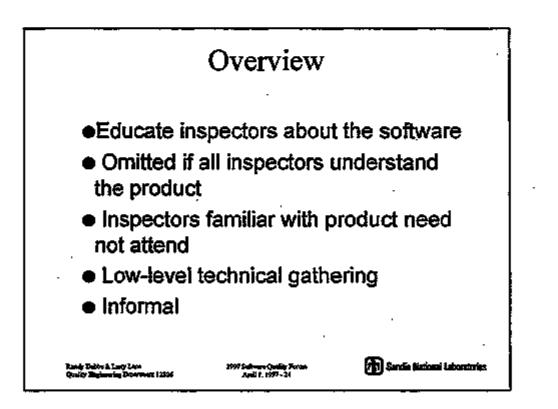
£

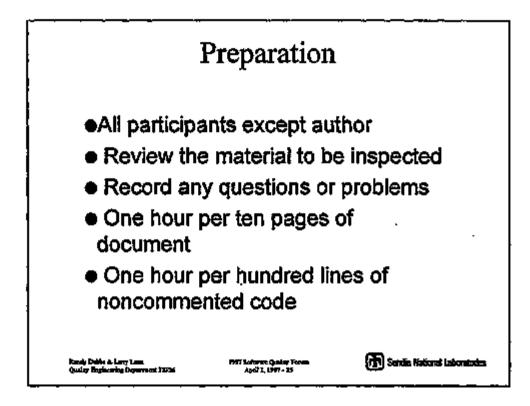


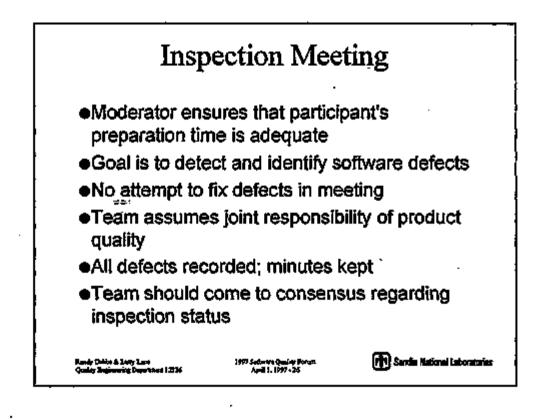


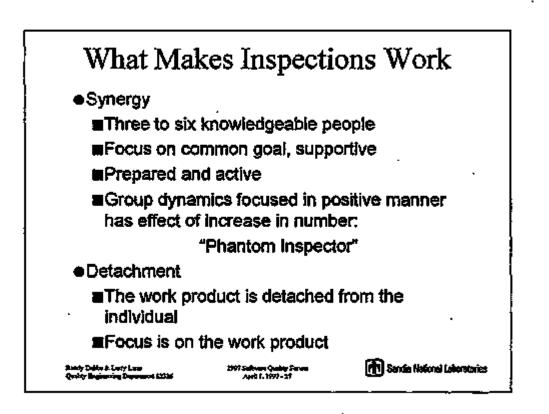
•

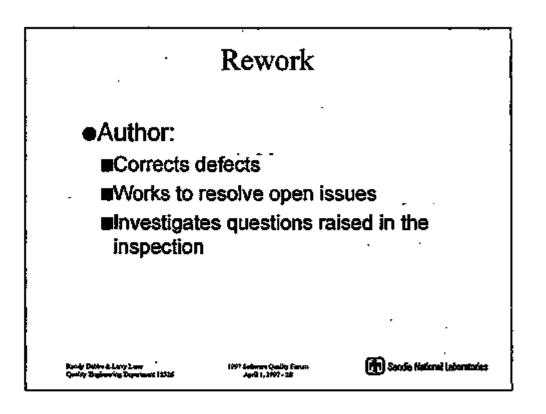


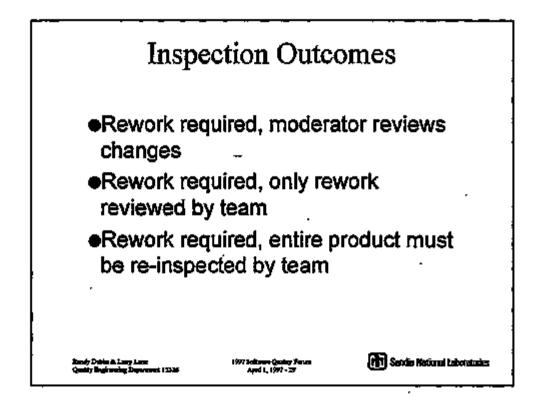


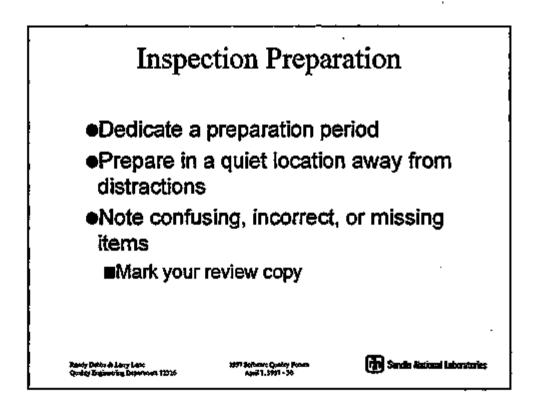


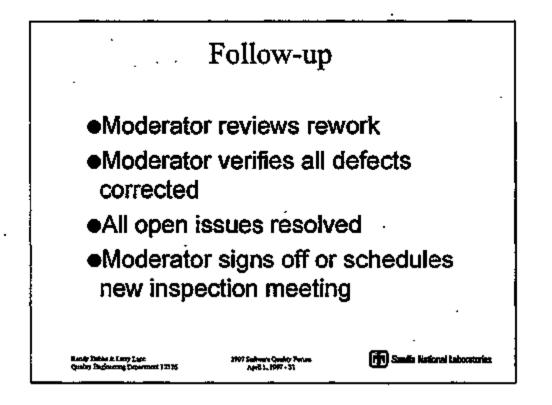


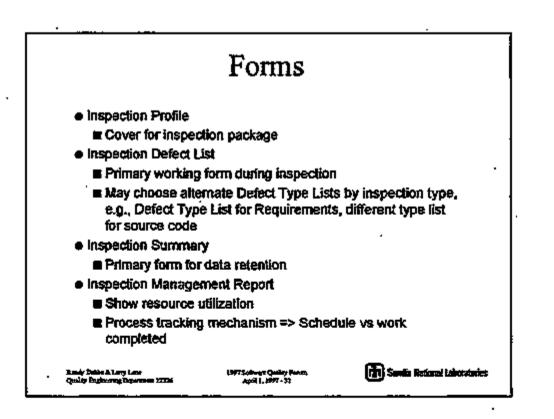




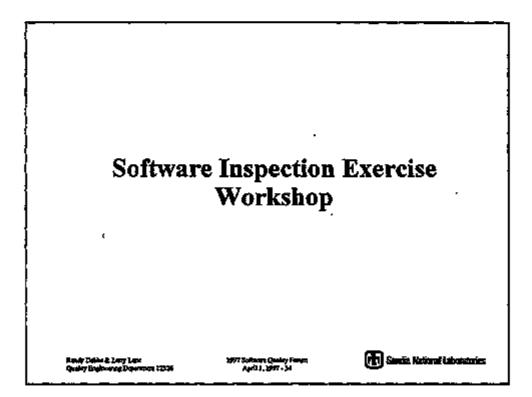


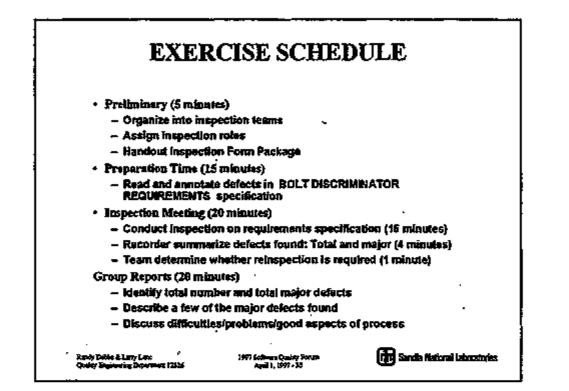


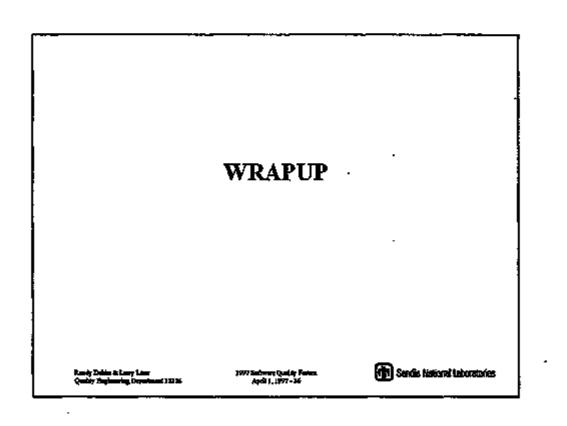




		inspectio	n Defect List				
Project	Document		P	etec	P	104: <u> </u>	·
inspection Type: 🖸 Requirements	O Oreign	0 000	D Test Pao	D Other			
Page Location Defect				Defect Type	Defect Severity	Defect Southe	Followii Chack
						<u> </u>	
		-	<u> </u>				<u> </u>
		· · ·					
							·
		_					
<u> </u>							
					<u> </u>		
		• • •					
Anthropology (Transmission)	Sergen La Martin La Martin		Cartan Gartan Sei Rassallan Sei Rassallan Sei Rass Sei Rass			Page Q	*. <u></u>
Randy Dables & Larry Larry		1977 Seitere	n Quality Forum (1997-11)	~⊗∧			



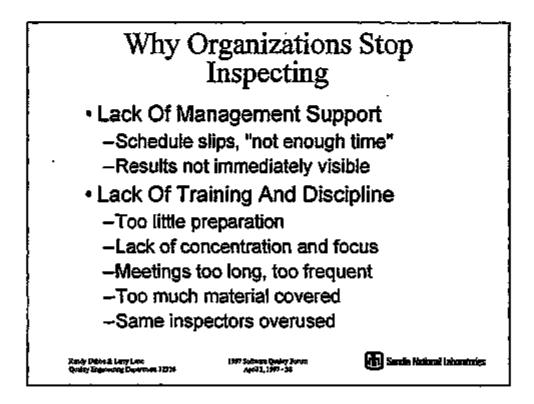




Guidelines for Successful Inspections

- Allow adequate preparation time
- Limit inspections to 2-hour sessions with no more than 2 sessions per day
- · identify problems; don't try to solve them
- Disassociate the author from the author's work
- Stress preparation, concentration and tolerance
- No management participation
- Choose the right participants

andy Dalais & Long Loon Juday Baginneray Dependent I 1997 Şalaraş Quelaş Fanası Ayrlı 1, 1997 - 27 🚯 Sanda National Laboratories



Recipe for Destroying an Inspection

- Invite your boss ...
- Invite everyone
- · Try to fix things
- · Make it last forever
- Do it on a Monday morning or Friday afternoon
- · Blitz through large amounts of material
- Get involved with personalities

Tooly Dable & Laty Late Quilty Businesing Date West 12526 1997 Salivite Quiley Polas April 1, 1997 - 39 👘 Sanda Vizional Vakoratning

Additional References

- "Experience with Inspection in Ultralarge-Scale Developments, " Russell, G. W., IEEE Software, January 1991, pp 25 - 31.
- "Getting Started on Metrics Jet Proplusion Laboratory Productivity and Quality," Bush, M. W., IEEE Experience Report, 1990.
- Structured Walkthroughs, Yourdon, Edward, Prentice - Hail, Englewood Cliffs, New Jersey, 1985.
- "Lessons from Three Years of Inspection Data," Weller, Edward, F., IEEE Software, September 1993, p.38 -45.
- "Annotated Bibliography on Software Inspections," Brykczynski, William, Software Engineering Notes, Vol., 18, No. 1, January 1993, pp 81-88.



Opening Session: Keynote Address

Capers Jones Chairman, Software Productivity Research Burlington, MA USA

Software Quality for 1997 - What Works and What Doesn't?

¢



Capers Jones

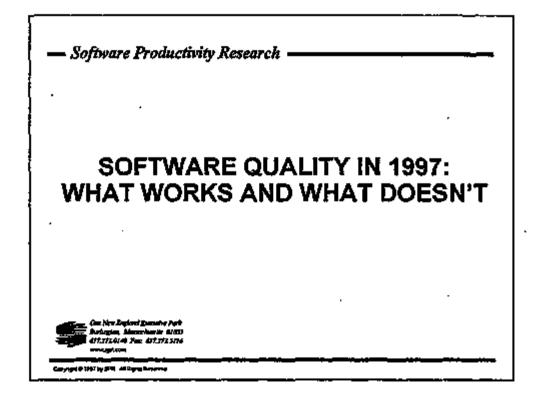
Keynote Address: Software Quality for 1997 - What Works and What Doesn't?

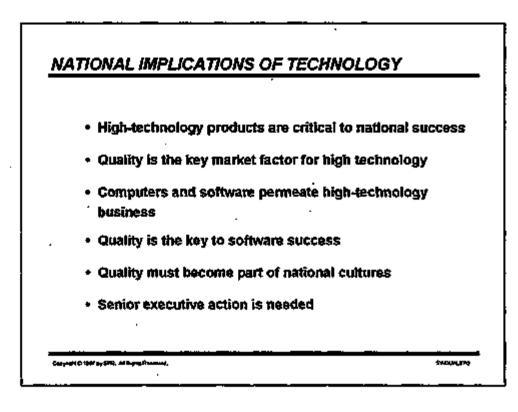
Capers Jones is an international consultant on software management topics and Chairman of Software Productivity Research, Inc. (SPR) in Burlington, MA. Following graduation from the University of Florida, Mr. Jones began his software career as a programmer in the office of the Surgeon General, Washington, D.C., Prior to becoming Chairman at SPR, Mr. Jones also worked at the Crane Company. IBM and was Assistant Director of Programming Technology at ITT in Stratford CT. Mr. Jones has published nine books dealing with software areas, including; Programming Productivity, Software Measurement, Software Quality. His tenth book, Software Cost Estimating is scheduled for publication in early 1997. Mr. Jones will share his experience and insights in his keynote address "Software Quality for 1997 -What Works and What Doesn't".

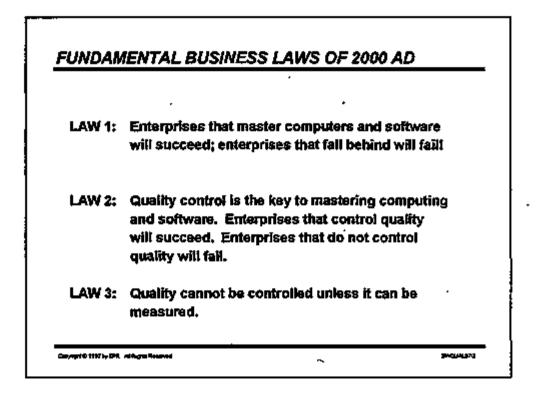
Keynote Address: April 2 1997, 09:00 - 10:00 am, TTC Auditorium

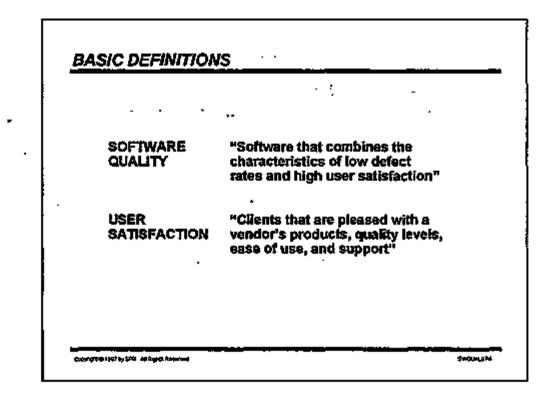
This presentation provides a view of software quality for 1997 – what works and what doesn't. For many years, software quality assurance lagged behind hardware quality assurance in terms of methods, metrics, and successful results. New approaches such as Quality Function Deployment (QFD) the ISO 9000-9004 standards, the SEI maturity levels, and Total Quality Management (TQM) are starting to attract wide attention, and in some cases to bring software quality levels up to a parity with manufacturing quality levels. Since software is on the critical path for many engineered products, and for internal business systems as well, the new approaches are starting to affect global competition and attract widespread international interest. It can be hypothesized that success in mastering software quality will be a key strategy for dominating global software markets in the 21st century.

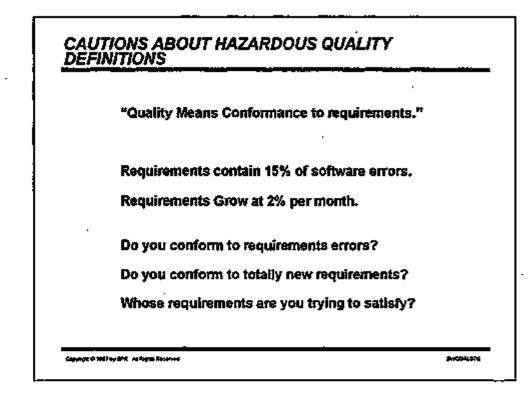
Capers Jones, Chairman Software Productivity Research, Inc. 1 New England Executive Park Burlington, MA 01803-5005 Phone 617 273 0140 FAX 617 273 5176 Email capers@spr.com

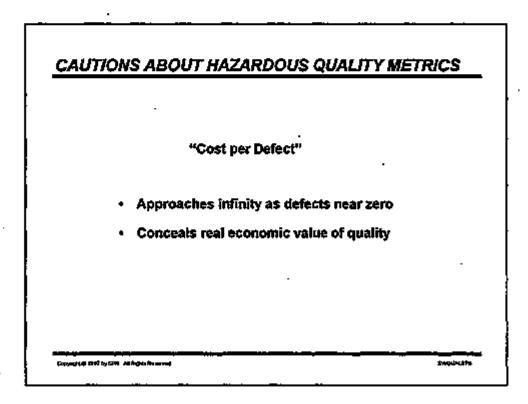






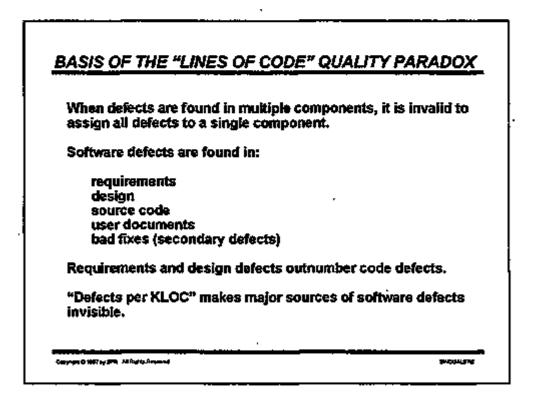




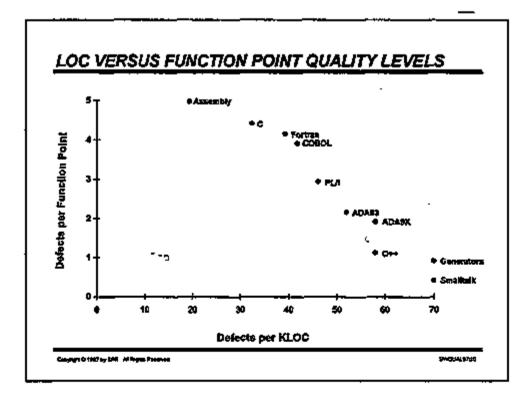


з

	(A) Poor Quality	(B) Good Quality	© Excellent Quality	D Zero Defects
Function Points	100	100	100	100
Bugs Discovered	500	50	5	0
Preparation	\$5,000	\$5,000	\$5,000	\$5,000
Removal	\$5,000	\$2,500	\$1,000	\$ 0
Repairs	\$25,000	<u>\$5,000</u>	<u>\$1.000</u>	<u>\$ 0</u>
Totel	\$35,000	\$12,500	\$7,000	\$5,000
Cost per Defect	\$70	\$250	\$1,400	\$\$
Cost per Function Point	\$350	\$125	\$70	\$50



Defect Origin	Assembly	<u>Ada</u>	<u>Objęctive Ç</u>	Full Rei
Requirements	35	35	35	15
Design	75	75	50	6
Code Im	165	25	10	2
Documents	50	50	50	10
Bad Fixes	25	15	5	2
TOTAL DEFECTS	300	200	150	35
Defects per KLOC	30	100	120	140
Defects per Function Point	6	4	2.4	0.7



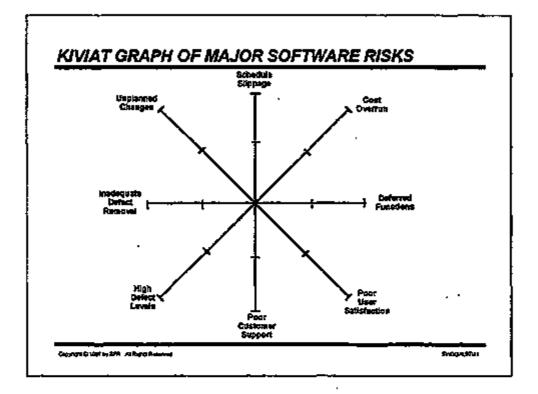
.

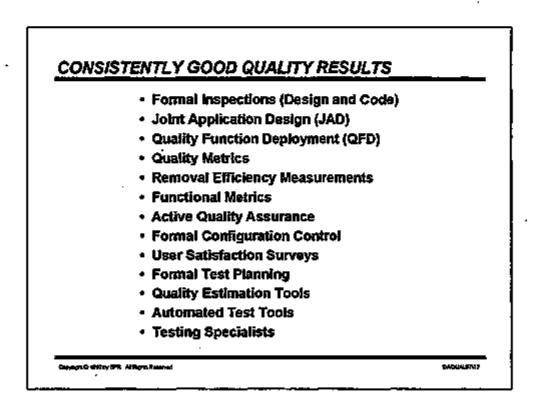
~

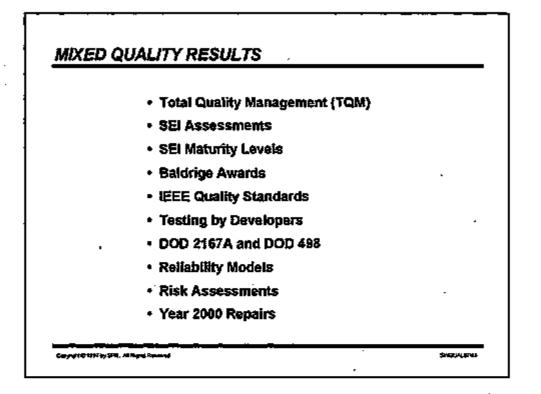
τ

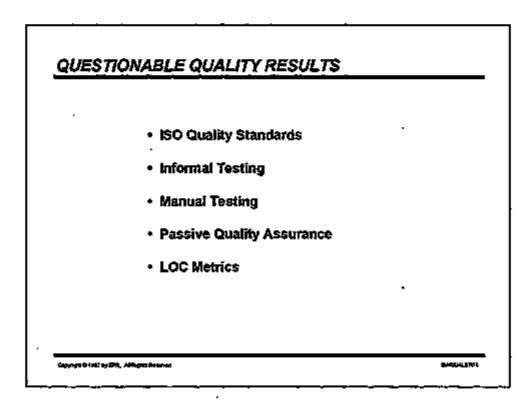
-

.









A PRACTICAL DEFINITION OF SOFTWARE QUALITY (PREDICTABLE AND MEASURABLE)

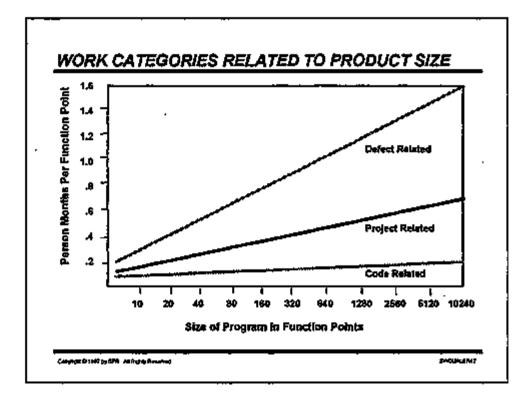
- Low Defect Potentials (< 2 per Function Point)
- High Defect Removal Efficiency (> 95%)
- Unambiguous, Stable Requirements (< 2.5% change)
- Explicit Requirements Achieved (> 97.5% achieved)
- High User Satisfaction Ratings (> 90% "excellent")
 - Installation
 - Ease of learning
 - Ease of use
 - Functionality
 - Compatibility
 - Error handling
 - User information (screens, manuals, tutorials)

Shikuk Mas

- Customer support
- Defect repairs

Country of the second

	SPR	<u>jso</u>
Defect Potential Estimation	Yes	Missing
Defect Removal Efficiency Estimation and Measurement	Yes	Missing
Delivered Defect Estimation and Measurement	Yes	Yes
User Satisfaction Measurement	Yes	Yes
Inspections and Reviews	Rigorous	Informal
Testing	Rigorous	Rigorous
Process Analysis	Rigorous	Informal



Size in <u>Function Points</u>	MgL/ <u>Support</u>	Defect <u>Remoyal</u>	<u>Paperwork</u>	<u>Coding</u>	<u>Totaj</u>
10,240	18%	35%	35%	12%	100%
5,120	17%	33%	32%	18%	100%
2,580	16%	31%	29%	24%	100%
1,260	15%	29%	26%	30%	100%
640	14%	27%	23%	36%	100%
320	13%	25%	20%	42%	100%
160	12%	23%	17%	48%	100%
80	11%	21%	14%	54%	100%
40	10%	19%	11%	60%	100%
20	9%	17%	8%	66%	100%
10	8%	15%	5%	72%	100%

	(D)	efects per Fun	ction Point))		
	System Software	Commercial Software	Information Software	Military Software	Overali Average
Defect Polentials	·6.0	5.0	4.5	7.0	5.6
Defect Remova) Efficiency	94%	94%	73%	96%	88%
Delivered Defects	0.4	0.5	1.2	0.3	9. 65
First Year Discovery Rate	65%	78%	30%	75%	60%
First Year Reported Defects	ŧ.25	0.35	0.36	0.23	0.30

U.S. SOFTWARE QUALITY AVERAGES

(Data Expresse	d in Terms of D	efects per Funci	tion Point)
<u>Defect Origins</u>	Defect <u>Potential</u>	Removal <u>Efficiency</u>	Delivered <u>Defects</u>
Requirements	1.00	77%	0.23
Design	1.25	85%	0.19
Coding	1.75	95%	0.09
Documents	0.60	80%	0.12
Bad Fixes	0.40	70%	0.12
TOTAL	5.00	85%	0.75
NCLUSIONS	¥6		
		85%	0.75

-)

(Data	a Expressed	t in terms of	Defects per	Function Poi	nt)
Size	Defect Potential	Defect Removal Efficiency	Delivered Defects	ist Year Discovery Rale	1st Year Reported Defects
1	1,85	95,00%	9.05	90.00%	0.08
10	2.45	92.00%	0.20	80.00%	0.16
108	3.68	90.00%	9.37	70.00%	0.26
1000	\$. 0D	\$5,00%	0.75	60.00%	0.38
10000	7.60	78.00%	1.87	40.00%	0.67
100000	8.65	75.00%	2.39	30.00%	0.72
VERAGE	5.02	86.53%	0.91	60.00%	0.38

.

.

SOFTWARE DEFECT REMOVAL EFFICIENCY AND THE FIVE LEVELS OF THE SEI CMM

· .

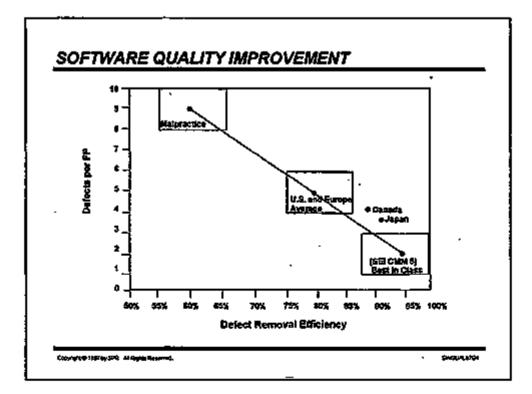
	Minimum	Average	Maximum
SEL Level 1	70.00%	65.00%	95.00%
El Level 2	70.00%	67.00%	96.00%
SEI Level 3	75.00%	89.00%	97,00%
SEI Level 4	80.00%	94.00%	99.00%
SEI Level 5	90.00%	97.00%	99.90%

- -

• .

ι,

(Data Expresse	d in Terms of Defec	ts per Function	Point)
SEI CMM Levels	Defect <u>Potentials</u>	Removal <u>Efficiency</u>	Delivered <u>Defects</u>
SEL CMM 1	⊊ [.] 5.00	85%	0.75
SEI CMM 2	4.00	90%	0.40
SEI CMM 3	3.00	95%	0.15
SELCIMM 4	2.00	97%	0.08
SELCIMM 5	1.00	89%	0.01



12

-

ر^

__)

ļ

ł

į

	Year 95% Exceeded (Approximate)
1. Telecommunications Manufacturing	1975
2. Computer Manufacturing	1977
3. Aero-space Manufacturing	1979
4. Military and Defense Manufacturing	1980
5. Medical Instrument Manufacturing	1980
6. Commercial Software Producers	1992



	Reliable Software	Software involved in Litigation for Poor Quality
Formal design Inspections	Used	Not used
Formal code inspections	Used	Not used
Subroutine testing	Used ,	Used
Unit testing	Used	Used
New function testing	Used	Rushed or omitted
Regression testing	Used	Rushed or omitted
Integration testing	Used	Used
System testing	Used	Rushed or omitted
Performance testing	Used	Rushed or omitted
Capacity testing	Úsed	Rushed or omitted

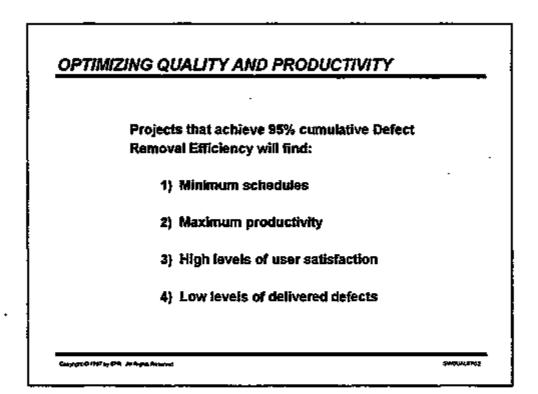
(Data Expre	ssed in '	Terms of	Defects Pe	er Functio	n Point)	
Year 1946	End-User	MIS	Outsrc.	Ĉommer.	System	Millary 1.59	Average 1.50
1950		2.00			2.50	2.00	2.17
1955		2.25			2.60	2.50	2.42
1960		2.59		1.50	3.00	3.00	2.50
1966		2.59		1.76	3.25	3.60	2.75
1970		2.75		2.59	4.00	4.50	3.44
1975	1.80	3.00	3.00	3.00	6.00	5,58	3.42
1980	1.50	3.76	3.60	3.50	6.00	6.25	4.68
1985	2.00	5.00	4.50	4.50	6.00	7.00	4.83
1890	2.50	5.00	4.75	4.75	6,50	7.00	5.08
1995	2.50	5.50	6.00	6,25	6,00	6.50	5,13
2000	3.00	6.00	5.50	6.00	6.50	6,50	5.58
Average	2.08	3.66	4,38	3.64	4.66	4.66	3.84

. -

(Data Ex)	pressed in T	erms of Pe	ercentage	of Defects F	Removed E	Sefore Dep	oloyment)
Year	End-User	Mis	Outerç.	Commer.	System	Military	Average
1546						\$0.00%	80.00%
1950		78.00%			83.00%	80.00%	80,33%
1965		79.00%	•		85,00%	86.00%	83,00%
1660		89.00%		80.00%	86.00%	86.00%	82,75%
1865		89.00%		\$2.00%	88.00%	86.00%	83.50%
1970		81.00%		84.00%	88.00%	\$7.00%	¥5.00%
1975	60.00%	\$2.00%	85.00%	\$5.00%	82.00%	90.00%	82.33X
1890	63.00%	82.00%	85.90%	89.00%	94.00%	91.00%	\$4.00%
1885	66.00%	84.00%	88.00%	90.00%	84,00%	92.00%	85,50%
1990	67.00%	84.00%	90.00%	92.00%	94.00%	83.00%	86,67%
1995	70.00%	85.00%	91.00%	84.00%	\$5.00%	95.00%	88,50%
2000	76,00%	88.00%	93.00%	95.00%	98.00%	96.00%	90.83%
Average	66,67%	82.09%	88.67%	87.89%	90.55%	88.33X	84,03%

ſ	Data Express	ed in Ter	ms of Defe	cts Deilvere	ed Per Fun	ction Polr	10
Year	End-Veer	MIS	Outsrc.	Commer.	System	Military	Average
1945					-	ē.30	0.30
1950		0.44			0.43	0.40	0.42
1956		0.47			0.38	0.38	0.41
1950		0.50		0.30	0.42	0.45	0.42
1965		0.50		0.32	0.46	0.49	0.44
1970		0.62		0.40	0.48	0.59	0.50
1975	0.40	0.54	0.45	0.45	0.40	0.55	0.47
1980	0.66	0.68	0.53	0.39	0.36	0.56	0.51
1985	0,70	0,80	0.54	0.45	0.38	0.58	0.57
1990	0.85	0.80	0.48	0.38	0.39	0.49	5.0 0.68
1995	0.75	0.83	0.46	0.32	0.24	0.33	0.48
2000	0.75	0.66	0.39	0.30	0.13	0.26	0.41
Average	0.66	0.61	0.AT	0.37	0.37	0.45	0.49

(Develop	(Development Defects + 1 Year of User Defect Reports)			
	PR <u>tance Level</u>	Efficiency Measured at One Year of Usage		
1. Exc	ellent	> 99%		
2. God	d	95%	•	
3. Ave	rage	87%		
4. Mar	ginal	83%		
5. Poc)Ľ	< 80%		



)

	ise defect remo origins is a val			ost element, studying
ІВМ С	orporation (MV	S)	SPR Cor	poration (client studies)
45%	Design error	\$	20%	Requirements errors
25%		s	30%	Design errors
20%	Bad fixes		35%	Coding errors
5%	 Documentati 	on errors	10%	Bad fixes
57	🛓 Administrativ	ve errors	<u> </u>	Documentation errors
100%	0		100%	
TRW C	orporation	Mitre Corp	oration	Nippon Electric Corp
60%	Design errors	64% De:	sign erron	5 60% Design errors
40%	Coding errors		ding error	
100%	-	100%	-	100%

Function points raised to the 0.3 power can predict the optimal number of defect removal stages.			
FUNCTION POINTS	DEFECT REMOVAL STAGES		
1	1		
10	2		
100	4		
1,000	8		
10,000	16		
100,000	32		
1,060,000	. 64		

Function points raised to the 1.2 power can predict the probable number of test cases for full test coverage.		
FUNCTION POINTS	TEST CASES	
1	1	
10	16	
100	251	
1,000	3,981	
10,000	. 63,096	
100,000	1,000,000	

RANGES OF TEST CASES PER FUNCTION POINT FOR SOFTWARE PROJECTS

Testing Stage	Min im um	Average	Maximum
Clean-room testing	0.60	1.00	3.00
Regression testing	0.49	0.60	5.30
Unit testing	0.20	0.46	1.29
New function testing	0.25	0.40	0.90
Integration testing	0.20	0.40	0.75
Subroutine testing	0,20	0.30	9.40
Independent testing	0.00	0.30	9.55
System testing	0.16	0.25	0.60
Viral texting	0.00	0.20	0.40
Petformance testing	0.00	0,20	0,40
Acceptance testing	0.00	0.20	0.60
Lab testing	60,0	0.20	0.58
Field (Beta) testing	0.00	0,20	1.08
Usability testing	0.00	0.20	0.40
Platform testing	0.00	0.15	0.30
Stress testing	0.00	0.15	0,30
Security testing	8.00	0.16	0.35
Year 2000 Testing	0.00	0.16	0,30
Total	2.00	5,50	13.26

18

`)

_)

Function points raised to th probable number of defects	e 1.25 power can predict the
(Defects in requirements, de bad fix categories.)	sign, code, documents, and
FUNCTION POINTS	POTENTIAL DEFECTS
1	1
10	18
\$00	316
1,000	5,623
10,000	100,000
100,000	1,778,279

	ΟΟυςτινιτγ
٠	The most effective way of improving software productivity and shortening project schedules is to reduce defect levels
•	Defect reduction can occur through:
	1. Defect prevention technologies
	Structured design
	Structured code
	High-level languages Etc.
	2. Defect <u>removal</u> technologies
	Design reviews
	Code inspections
	Tests
	Correctness proofs

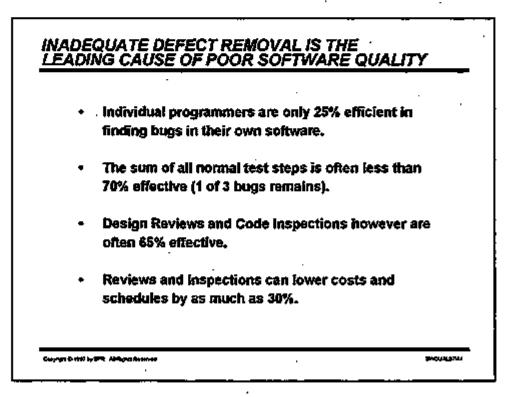
	Requirements Defects	Design Defects	Code Defects	Document Defecte	Performance Defects
JAD's	Excelient	Good	Not Applicable	Fair	Poor
Prototypes	Excellent	Excellent	· Fair	Not Applicable	Excellent
Structured Methods	Fair	Good	Excellent	Fair	Fair
CASE Tools	Fair	Good	Fair	Fair	Fair
Blueprints & Reusable Code	Excellent	Excellent	Excellent	Excellent	Good
QFD	Good	Excellent	Fair	Poor	Good

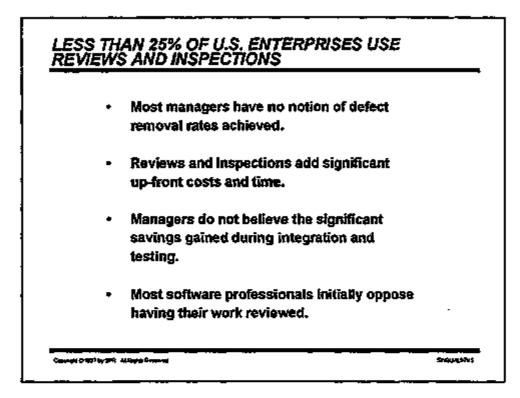
	Requirements Defects	Design Defects	Code Defects	Document Defects	Performance Defects
Reviews/ Inspections	Faît	Excellent	Excellent	Good	Fair
Prototypes	Good	Fair	Feir	Not Applicable	Good
Testing (all forms)	. Poor	Poor	Good	· Fair	Excellent
Correctness Proofs	Poor	Poor	Good	Fair	Poor

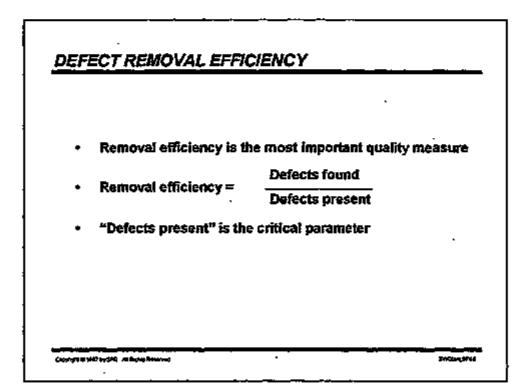
	Melhods	Training	Experience	Enthusiaam	Nanagemen Support
1. Excellent	Format	Formal	Substantial	Good	Good
2. Good	Formal	Formal	Mixed	Good	Moderate
3. Average	Informal	Informat	Mixed	Mixed	Mixed
4. Marginal	Informal	Informal	Little	Minimat	Minimal
5. Poor	informal	Informal	None	Negative	Minimal

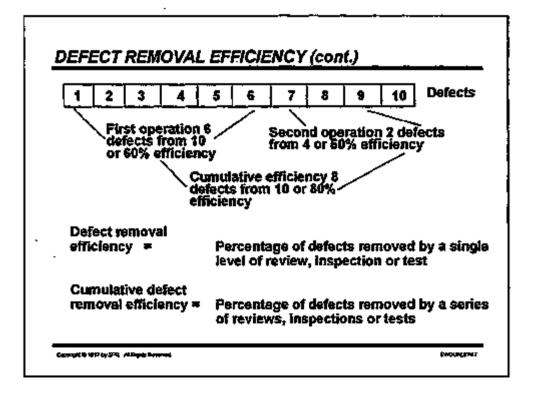
				•				
		Defect Estimation	Defect Tracking		Complexity <u>Measures</u>	Test Coverage <u>Measures</u>	Remôvaŭ <u>Measures</u>	Maintenan <u>Measure</u> t
1.	Excellent	Yes	Yes	Yes	Yes	Yes	Yes	Yes
2	Good	Yes	Yes	Yes	No	Yes	No	Yes
3.	Average	No	Yea	Yes	No	Yes	No	Yes
4.	Marginal	No	No	Yes	No	Yes	No	Yes
6.	Poor	No	No	No	No	No	No	No

(Tool Capacity Expressed in Function Points)						
Fool Categories	Lagging	Average	Leading			
Statistical analysis tools			3,00D			
Quality estimation models			2,500			
Spreadsheet	750	1,250	2,000			
Graphica/Presentations	750	1,250	2.000			
Word processing	500 500	1,000 1,250	2,000			
Configuration control	500	1,250	2,000			
Test case generators			1.750			
Data base	500	1,000	1,500			
Defect tracking/Analysis	500	750	1,000			
Reliability estimation models		500	1,000			
Symbolic debuggers	250	500 500	750			
Electronic mail	300	500	700			
Appointment calendar	100	300	750			
Phone/Address file	100	150	500			
Complexity analyzers			350			
Test path coverage analyzers		200	· 350			
Test execution monitors		200	350			
Totals	4,250	8,850	22,250			









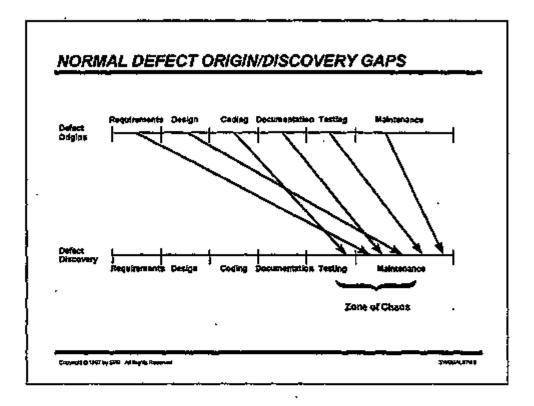
	Lowest	<u>Median</u>	<u>Hiahest</u>
Requirements review	20%	30%	50%
Top-level design reviews	30%	40%	60%
Detailed functional design reviews	30%	45%	65%
Detailed logic design reviews	35%	55%	75%
Code inspections	35%	60%	85%
Unit tests	10%	25%	50%
Function tests	20%	35%	55%
Integration tests	25%	45%	60%
Site/installation tests	<u>25%</u>	<u>50%</u>	<u>65%</u>
	75%	95%	99%

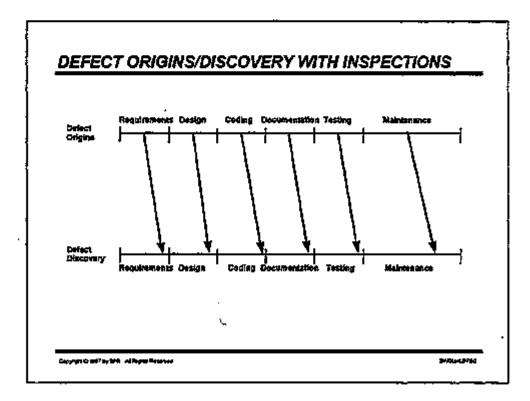
CT 7.4-CL

24

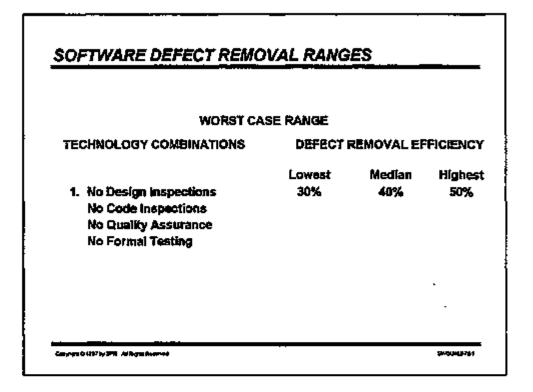
12.

L





.



SINGLE, TECHNOLOGY CHANGES ECHNOLOGY COMBINATIONS DEFECT REMOVAL EFFICIENCY			
2. No design inspections No code inspections FORMAL QUALITY ASSUR/ No formal testing	Lowest 32% NICE	Median 45%	Highest \$5%
3. No design inspections No code inspections No quality assurance FORMAL TESTING	. 37%	53%	60%
4. No design inspections FORMAL CODE INSPECTIO No quality assurance No formal testing	43% NS	57%	65%
 FORMAL DESIGN INSPECT No code inspections No quality assurance No formal festing 	ions 45%	60%	68%

• •

÷

ì

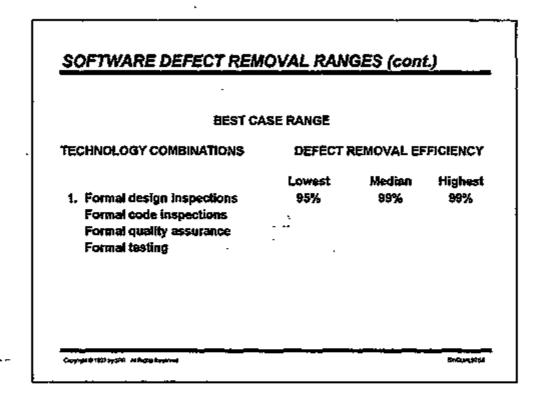
-)

TWO TECHNOLOGY CHANGES				
TECHNOLOGY COMBINATIONS	DEFECT REMOVAL EFFICIE		DEFECT RE	FICHENCY
	Lowest	Median	Higbest	
6. No design inspections No code inspections FORMAL QUALITY ASSURANCE FORMAL TESTING	50%	65%	75%	
7. No design inspections FORMAL CODE INSPECTIONS FORMAL QUALITY ASSURANCE No formal testing	53%	68%	78%	
8. No design inspections FORMAL CODE INSPECTIONS No quality assurance FORMAL TESTING	55%	70%	60%	

TWO TECHNOLOGY	CHANGES	(cont.)	
TECHNOLOGY COMBINATIONS	DEFECT	REMOVAL EP	FICIENCY
9. FORMAL DESIGN INSPECTIONS No code inspections FORMAL QUALITY ASSURANCE No formal testing	Lowest 60%	Median 75%	Highest 85%
10. FORMAL DESIGN INSPECTIONS No code inspections No quality assurance FORMAL TESTING	55%	80%	87%
11. FORMAL DESIGN INSPECTIONS FORMAL CODE INSPECTIONS No quality assurance No formal testing	70%	86%	90%

٠,

	THREE TECHN	DLOGY CHANG	ES 🛛	
TEC	HNOLOGY COMBINATIONS	DEFECT REMO	VAL EFFICIEN	CY
		Lowest	Median	Highest
12.	No design inspections FORMAL CODE INSPECTIONS FORMAL QUALITY ASSURANCE FORMAL TESTING	75%	87%	83%
13.	FORMAL DESIGN INSPECTIONS No code inspections FORMAL QUALITY ASSURANCE FORMAL TESTING	77%	90%	85%
14.	FORMAL DESIGN INSPECTIONS FORMAL CODE INSPECTIONS FORMAL QUALITY ASSURANCE No formal testing	85%	35%	97%
15.	FORMAL DESIGN INSPECTIONS FORMAL CODE INSPECTIONS No quality desurance FORMAL TESTING	85%	97%	99%



ないりんて

Defect Removal Efficiency Level (Percent)	Number of Projects	Percent of Projects
> 99	6	0.40%
95 - 99	104	6.93%
90 - 95	263	17.53%
85 - 90	559	37.26%
80 - 85	406	27.20%
< 80	161	10.73%
Total	f,500	100.00%

DISTRIBUTION OF 1500 SOFTWARE PROJECTS BY DEFECT REMOVAL EFFICIENCY LEVEL

APPROXIMATE DISTRIBUTION OF TESTING METHODS FOR U.S. SOFTWARE PROJECTS

•

Testing Stage	Percent of Projects Utilizing Test Stage
General Forms of Testing	
Subroutine testing	100%
Unit testing	99%
System testing of full application	95%
New function testing	90%
Regression testing	70%
Integration testing	.50%
Specialized Forms of Testing	
Viral protection testing	45%
Stress or capacity testing	35%
Performance testing	30%
Security testing	15%
Platform testing	5%
Year 2000 testing	5%
independent testing	3%
orgen und in site setting diversion	

Testing Stage	Percent of Projects Utilizing Test Stage
Forms of Testing Involving Users	
Customer acceptance testing Field (Beta) testing Usability testing Lab testing Clean-room statistical testing	35% 30% 20% 1% 1%

(Size of Application in Function Points)							
Class of Softwar	1	10	100	° 1K	1 0K	160K	Average
End-uzer	1	2	2				1.57
MIS	2	3	4	6	7	8	5.00
Outsourcers	2	3	5	7	8	9	5.87
Commercial	3	- 4	6	9	11	12	7.50
Systems	3	4	7	11	12	14	B.50
Military	4	6	8	11	13	16	8,50
Average	2.5D	3.60	5.33	8.8D	10.20	11.80	7.02

NUMBER OF TESTING STAGES, TESTING EFFORT, AND DEFECT REMOVAL EFFICIENCY

Devoted to Testing	Removal Efficiency
10%	50%
15%	60%
20%	70%
25%	75%
30%	80%
33%*	85%*
36%	87%
39%	90%
42%	92%
, 33% costs, and 85% removal	efficiency are U.S. averages.
	15% 20% 25% 30% 33%* 36% 39% 42%

NUMBER OF TESTING STAGES, TESTING EFFORT, AND DEFECT REMOVAL EFFICIENCY (cont.)

Number of	Percent of Effort	Cumulative Defect
Festing Stages	Devoted to Testing	Removal Efficiency
10 testing stages	45%	94%
1 testing stages	48%	° 95%
12 testing stages	52%	98%
13 testing stages	55%	99%
14 testing stages	58%	99.9%
15 testing stages	61%	99.99%
6 testing stages	64%	99.999%
17 testing stages	67%	99,9999%
8 testing stages	70%	99,99999%
Note: Six test stage:	s, 33% costs, and 85% removal	efficiency are U.S. averages.
		- •
Canada D 1987 to STR. Al Rota farmer		\$MOUK376

-

٠

J,



- No single method is adequate.
- Testing alone is insufficient.
- Reviews, inspections and tests combined give high efficiency, lowest costs and shortest schedules.
- Reviews, inspections, tests and prototypes give highest cumulative efficiency.
- Administrative problems need special solutions.
 Ordinary defect removal is not adequate.
- Maintenance costs are cumulative, expensive and chronic.

Conset @ 1987 by SPR 48-0 ptg General

Software Quality in 1997

January 9, 1997

Abstract

For many years, software quality assurance lagged behind hardware quality assurance in terms of methods, metrics, and successful results. New approaches such as Quality Function Deployment (QFD) the ISO 9000-9004 standards, the SEI maturity levels, and Total Quality Management (TQM) are starting to attract wide attention, and in some cases to bring software quality levels up to a parity with manufacturing quality levels. Since software is on the critical path for many engineered products, and for internal business systems as well, the new approaches are starting to affect global competition and attract widespread international interest. It can be hypothesized that success in mastering software quality will be a key strategy for dominating global software markets in the 21st century.

Capers Jones, Chairman Software Productivity Research, Inc. 1 New England Executive Park Burlington, MA 01803-5005

з

Phone 617 273 0140 FAX 617 273 5176 Email capers@spr.com

> Copyright © 1995-1997 by Capers Jones, Chairman, SPR, Inc. All rights reserved.

INTRODUCTION

Software has become one of the most pervasive technologies of the 20th century. Within the past 30 years, software has spread from a small number of comparatively specialized applications to become a critical factor in almost all engineered products. Software has also become a major factor in consumer goods, and in company operations, military operations, and government operations. Thirty years ago, poor software quality was often annoying, but today poor software quality can literally shut down a phone system, a defense system, and even a company. Any reasonable prognosis makes software even more critical in the future, and hence software quality will become more critical than today as well.

As many countries strive to compete in the international software market place, quality is now a major topic for both software vendors and for outsource contractors. Any country or company that wants to achieve a major place in world software markets must achieve and maintain high software quality levels.

Barriers to Software Quality Exploration

Progress in all forms of engineering is heavily dependent upon accurate measurement and precise metrics. Software achieved notoriety as being the worst measured engineering discipline of all time. The main barrier to software quality control in the 1950's, 60's, and 70's was a simple lack of good quantitative data about software quality levels, reliability, defect removal efficiency and other basic quality data. This lack of data was not because software managers and professionals did not care about quality, but because there were no effective metrics prior to 1979 that could actually be used to measure software quality.

Historically, software quality was measured crudely in terms of "defects found per 1000 source code statements" (normally abbreviated to KLOC). Unfortunately, that metric contained a built-in paradox which caused it to give erroneous results when used with newer and more powerful programming languages, such as Ada, object-oriented languages, or program generators. The results were so poor that several leading companies stopped trying to measure software, and lagging companies never started.

In 1979, A.J. Albrecht of IBM published a new metric for measuring both software quality and productivity, which he termed "Function Points." A Function Point is a synthetic metric derived from five visible external characteristics of software applications: 1) Inputs; 2) Outputs; 3) Inquiries; 4) Logical files; 5) Interfaces.

Function Points are completely divorced from lines of source code. In a sense, Function Points are like European Currency Units (ECU), which are synthetic metrics that allow rational economic and financial studies across multiple national currencies. Function Points allow rational quality and productivity studies across the 400 or so programming languages that have come into being.

In 1986, Function Point users formed a non-profit association, the International Function Point Users Group, or IFPUG. This organization and its affiliates now have over 500 corporations and government agencies as members in the United States, Canada, Europe, South America, and the Pacific Rim and membership is growing by more than 45% per year.

It is an interesting business phenomenon that measurement of software quality and productivity is now among the most rapidly growing technologies in the entire history of software.

One of the advantages of the Function Point metric is that it can be used to predict and measure all sources of software errors, and not just coding errors. Based on a study of more than 6700 software projects published in the book <u>Applied Software Measurement</u> (McGraw-Hill, 1996), the average number of software errors is about five per function point, apportioned across the following major defect origins. However, the "best in class" software organizations are achieving defect potentials of roughly half the total of "average" groups as shown in Table 1:

Defect Origins	Average Defects per Function Point	Best in Class Defects per Function Point	Difference
Requirements	1.00	0.40	0.60
Design	1.25	0.60	0.65
Coding	1.75	1.00	0.75
Document	0.60	0.40	0.20
Bad Fixes	0.40	0.10	0.30
Total	5.00	2.50	2.50

Table 1: U.S. Average	es and "Best in (Class" Defects	per Function Point
-----------------------	-------------------	----------------	--------------------

These numbers represent the total numbers of defects that are found and measured from early software requirements throughout the remainder of the lifecycle of the software.

Complementing the Function Point metric are measurements of defect removal efficiency, or the percentages of software defects removed prior to delivery of the software to clients. The U.S. average for defect removal efficiency, unfortunately, is currently only about 85% although the best projects in leading companies such as Motorola, Raytheon, IBM, and Hewlett Packard achieve defect removal efficiency levels well in excess of 99%.

All software defects are not equally easy to remove. Requirements errors, design problems, and "bad fixes" tend to be the most difficult. Thus, on the day when software is actually put into production, the average quantity of latent errors or defects tends to be about 0.75 per Function Point, with the following distribution as shown in Table 2:

Defect Origins	Defect Potentials	Removal Efficiency	Delivered Defects
Requirements	1.00	77%	0.23
Design	1.25	85%	0.19
Coding	1.75	95%	0.09
Document	0.60	80%	0.12
Bad Fixes	0.40	70%	0.12
Total	5.00	85%	0.75

Table 2: U.S. Averages for Defect Potentials and Removal Efficiency Levels

The best companies are using state-of-the art methods to lower their defect potentials, and coupling that with state-of-the-art methods for removing defects with high efficiency in excess of 95%. The results can be quite impressive.

COMPARING U.S. QUALITY DATA WITH INTERNATIONAL DATA

The author's company, Software Productivity Research, collects data on both productivity and quality in more than 20 countries. Although that may sound like quite a lot, it is still only a small and partial step toward a true global survey of software quality.

From the data collected, provisional averages on international quality levels were published in 1993 in the author's book, <u>Software Productivity and Quality Today - The</u> <u>Worldwide Perspective</u> (Information Systems Management Group, Carlsbad, CA). Following are excerpts from some of the preliminary global findings, with some data revised during 1995 and 1996:

Country	Defect Potential per Function Point	Defect Removal Efficiency Levels	Delivered Defects per Function Point
Japan	4.50	93%	0.32
Canada	4.55	86%	0.64
United States	5.00	85%	0.75
Norway	4.95	84%	0.79
Sweden	5.00	84%	0.80
France	4.75	83%	0.82
Italy	4.85	83%	0.82
India	5.10	84%	0.82

Table 3: International Comparisons of Defect Potentials and Defect Removal

Germany	4.95	83%	0.84
England	4.85	82%	0.87
South Korea	5.20	83%	0.88
Russia	5.50	80%	1.10

The margin of error of this data is very high, except for the United States, and the information is presented primarily to generate discussion about the two key topics of defect potentials and defect removal efficiency levels.

Within every country where the author and his colleagues have collected data, the ranges of defect potentials and removal efficiencies are very broad. Some companies are achieving potentials of less than 2 defects per function point and eliminating more than 95%, while other companies have defect potentials approaching 10 per function point and eliminate barely 75%.

Although for every country, the range of performance is quite broad some six industries stand out internationally as achieving the best overall software quality levels:

Industries With Best Software Quality Results

- 1. Computer manufacturers
- 2. Telecommunication equipment manufacturers
- 3. Defense and weapons system manufacturers
- 4. Aerospace manufacturers
- 5. Medical equipment manufacturers
- 6. Commercial software manufacturers

Companies within these six industries typically average more that 95% in cumulative defect removal efficiency, which places them well above the norms of the 40 industries for which SPR has collected quality data.

Four characteristics set these industries apart from industries with less effective quality control approaches: 1) Usage of formal design and code inspections; 2) Usage of formal and active quality assurance functions; 3) Usage of trained testing specialists and formal testing departments; 4) Usage of a powerful suite of defect estimation, defect tracking, and other quality control tools.

A common characteristic of these industries in every country is that much of their software controls physical devices such as computers, switching systems, weapons systems, aircraft, and the like. The single exception is that of the commercial software vendors, and in this industry it has been learned by trial and error that poor quality loses business.

TOOLS AND METHODS USED BY BEST IN CLASS QUALITY PRODUCERS

There are major variances from company to company and country to country in the sets of tools and methodologies used to approach software quality. However, the best in class organizations have a common nucleus which includes these factors:

Quality Measurements

The most striking difference between leading organizations and lagging ones in every country is that, without exception, the leaders know their quality levels and user satisfaction levels because they measure these factors very carefully.

The quality measurements in leading companies vary slightly, but usually include these elements: 1) Software defect volumes are measured from requirements or design throughout the rest of the development cycle and into the field; 2) Defect severity levels are measured, ranging from serious through minor; 3) Defect origins are measured, so that problems with requirements, design, code, documents, and secondary problems are known.

This software quality data is collected on a daily basis, and then summarized at monthly, quarterly, and annual intervals to show trends over time. In addition, the leaders also measure user satisfaction, although the frequency of user surveys is normally once or twice a year.

Quality Methods

The leading companies did not become good overnight. Most of them have been engaged in software quality control work for 20 years or more. Therefore the leading companies have developed a set of proven methods that are known to work. These methods are sometimes defined under two headings, *defect prevention* and *defect removal*. Here are some examples: 1) Formal inspections of design, code, and other deliverables are used by essentially all software quality leaders since these activities are highly effective in both preventing and removing software defects; 2) Active and energetic software quality assurance groups, which may exceed 5% of total staff, are often found in the industry leaders.

A very interesting correlation is that in every country the best in class quality producers tend to utilize formal inspections of design, code, and other deliverables. Formal inspections are one of the few kinds of defect removal operation to exceed 60% in defect removal efficiency, and on average are about twice as efficient as any common form of testing. (High-volume external Beta testing by more than 1000 clients simultaneously is the only form of testing that is more efficient in defect removal than inspections.)

Both industry leaders and laggards test their software. The most striking difference between leaders and laggards is what the leaders do before testing begins. By means of

defect prevention approaches such as Joint Application Design (JAD), Quality Function Deployment (QFD), formal inspections, and various flavors of structured analysis and design, the leaders usually have far fewer problems attributable to the front of their software development life cycles. Therefore when testing begins, the code developed by the leaders is substantially free from serious problems long before testing even starts. This translates into quicker testing cycles and fewer delays of final delivery.

Two important topics do not yet have any strong empirical correlations with software quality results: ISO 9000-9004 certification and the SEI capability maturity levels. Although the ISO standards are aimed at quality, they have not yet created any significant results within the software industry.

Indeed, as this report is being drafted a world wide web conference is on-going, hosted by John Seddon of the United Kingdom, to discuss whether or not ISO certification *degrades* quality rather than enhances it. In late 1996 a British "watch dog" government agency directed the British Standards Institute to stop making claims that ISO certification improved productivity or quality without empirical evidence to support the claims.

The SEI maturity level concept is also surprisingly ambiguous in terms of quality. There is a lot of overlap among the various SEI levels, and a surprising observation is that the worst software that is created by SEI level 3 organizations in terms of quality can lag the best software created by level 1 organizations.

However, some recent studies within the past two years do indicate an overall improvement in quality as SEI levels climb upward from level 1 to 3, 4, and 5. Unfortunately, the total number of samples is too small for statistical certainty.

Following are the current ranges of software defect potentials and removal efficiency levels observed from among client organizations that have utilized the SEI CMM:

Level 1 Quality: The software defect potentials noted from several hundred projects in Level 1 organizations run from about 3 to more than 15 defects per function point but average about 5.0 defects per function point. Defect removal efficiency runs from less than 70% to more than 95% but only averages about 85%. Thus the average number of delivered defects for Level 1 organizations is about 0.75 defects per function point.

Level 2 Quality: The software defect potentials noted from about 50 projects in Level 2 organizations run from about 3 to more than 12 defects per function point but average about 4.8 defects per function point. Defect removal efficiency runs from less than 70% to more than 96% but averages about 87%. Thus the average number of delivered defects for Level 2 organizations is about 0.6 defects per function point.

Level 3 Quality: The software defect potentials noted from about 30 projects in Level 3 organizations run from about 2.5 to more than 9 defects per function point but average about 4.3 defects per function point. Defect removal efficiency runs from less than 75%

to more than 97% but averages about 89%. Thus the average number of delivered defects for Level 3 organizations is about 0.47 defects per function point.

Level 4 Quality: The software defect potentials noted from 9 projects in Level 4 organizations run from about 2.3 to more than 6 defects per function point but average about 3.8 defects per function point. Defect removal efficiency runs from less than 80% to more than 99% but averages about 94%. Thus the average number of delivered defects for Level 4 organizations is about 0.2 defects per function point.

Level 5 Quality: The software defect potentials noted from 4 projects in a Level 5 organization ran from about 2 to 5 defects per function point but currently seem to average 3.5 defects per function point. Defect removal efficiency ran from less than 90% to more than 99% but averaged about 97%. Thus the average number of delivered defects for a Level 5 organization is about 0.1 defects per function point although there is obviously an insufficient sample at this level.

To illustrate the overlap of quality among the five levels of the SEI CMM, the following table shows our minimum, average, and maximum numbers of delivered defects per function point for each of the five CMM levels. Note that the best results from Level 1 are actually better than the worst results from Levels 3 and 4, even though the average results improve as the CMM ladder is climbed.

Table 4: Software Delivered Defects at Each Level of the SEI CMM

(Defects expressed in terms of defects per function point)

Minimum Average Maximum

SEI Level 1	0.150	0.750	4.500
SEI Level 2	0.120	0.624	3.600
SEI Level 3	0.075	0.473	2.250
SEI Level 4	0.023	0.228	1.200
SEI Level 5	0.002	0.105	0.500

Although samples are small for the higher levels, there is now evidence from studies such as the ones carried out by Software Productivity Research (SPR) in 1994 which indicate that when organizations do move from CMM level 1 up to the higher levels their productivity and quality levels tend to improve, although there is quite a bit of overlap among the five CMM stages.

Quality Tools

What is easily the most visible difference between industry quality leaders and quality laggards is the set of tools available to the leaders, and totally absent from the lagging organizations. The leaders usually employ a set of quality tools that include some or all of the following: 1) Quality estimation predictive tools; 2) Defect and quality measurement tools; 3) Test planning tools; 4) Test coverage analysis tools; 5) Software

reliability predictive models; 6) Complexity analysis tools; 7) Statistical analysis and reporting tools.

These tools have the general characteristic of putting quality in tangible, quantitative terms so that the underlying root causes can be explored and improved. The laggards tend to have no quantitative data, and hence are unable to take any kind of carefully planned corrective actions.

Since each of the quality tools cited in this section is roughly 1000 function points in size, it can be asserted that the leading quality assurance groups have in the range of 6000 to 8000 function points of quality-related tools available. By contrast, laggards with marginal quality levels often have less than 500 function points of quality-related tools, or even none at all.

Quality Culture

A final aspect which separates the laggards from the leaders is the culture of quality among the leaders, and its absence among the laggards. The word "culture" does not have a very precise definition, so in this context the meaning is the following: when visiting the industry leaders, almost everyone you talk to cares about quality and many of them also know something about it. When visiting the laggards, you tend to find some people who care about quality of course, and a few people who know how it might be achieved, but these quality-conscious people often feel isolated and even angry that their executives have no particular interest in the subject. There is no substitute for executive awareness of the importance of quality. When you meet an executive vice president or a CEO that can carry on a serious conversation about software quality, you can be fairly sure that the company is a pretty good one. When you visit a company where the executives know nothing of quality and give the appearance of not caring either, you can be fairly sure that the company will have some tough times ahead.

SUMMARY AND CONCLUSIONS

Now that the dimensions of software quality can be measured, it is obvious that there are two powerful sets of technologies which must both be deployed in order to be successful with software: 1) Defect prevention methods; 2) Enhanced defect removal methods.

The set of defect prevention methods includes all technologies which can simplify complexity, and minimize the tendency to make errors. Examples of software defect prevention methods include Joint Application Design (JAD), prototyping, structured methods, clean-room development, Information Engineering (IE), Object-Oriented methods (OO), Quality Function Deployment (QFD), and of course software quality measurement programs. Synergistic combinations of defect prevention methods can reduce defect potentials by more than 50% across the board, with the most notable improvements being in some of the most difficult problems, such as requirements errors. The set of defect removal methods include structured walkthroughs, formal inspections, audits, independent verification and validation, and many forms of testing. Accurate measurement of defect removal efficiency has revealed some surprising findings. One surprise is that most forms of testing are less than 30% efficient in actually finding software problems, due in part to the fact that test cases are almost worthless for finding requirements errors, and not terribly effective in finding design errors. Against front-end requirements and design defects, formal inspections often achieve more than 60% defect removal efficiency rates.

The "best in class" software producers now have defect potentials of less than 2.0 errors per Function Point, coupled with defect removal efficiencies that hover around 99% and may exceed it for mission-critical software. This combination yields delivered defect totals of only 0.02 defects per Function Point, or more than an order of magnitude better than U.S. norms and provisional international norms as well.

It can be hypothesized that international competition in the software domain will intensify as we move to the end of the 20th century. Since high levels of software quality are associated with high market shares, quality control is now a major topic of global competitiveness.



Session A1: Software Management

Chair Don Schilling AS/FM&T

Session : Paper #	Author(s)	Title
<u> </u>	Rodema Ashby	The Right Rock: Finding/Refining
	Sandia National Laboratories	Customer Expectations
A1:2	David Harris	TCAMS Lessons Learned
	Sandia National Laboratories	
A1:3	Joe Schofield	The Next Silver Bullet - Or Just
	Sandia National Laboratories	Another Shot in the Foot?

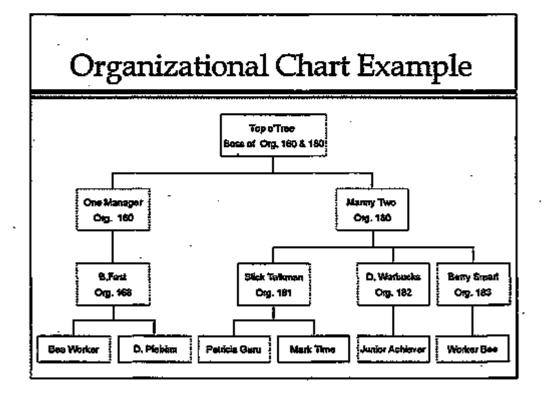
The Right Rock:		
Finding & Refining Customer		
	Expectations	
Finding:	Organization Chart Review Customer Interviews Customer Desires Matrix	
Refining:	Quality Functional Deployment Child Design Matrix Requirements Document Acceptance Test Document Create the User Manual Rapidly Prototype if Configurable Incrementally Build if Custom Dev.	
Roder	na Ashby, 844-2087, mrashby@sandia.gov Seda interdimpentationary genety Seda Copersion - Laboration Corpusy. Set to University Transmission - Company on the Company Device of Company on the Company Device of Company on the Company of the Company of the Company on the Company of the Company	

Can't Tell the Players Without a Scorecard

- Who is the Customer?
 - » The person using the system?
 - » Your Manager? Other's Managers?
 - » The person who paid for the development?
 - » A Sandia Initiative Director?
- A Stakeholder is anyone who will assess & affect the project success
 - » You don't get to pick, & ignorance is not bliss

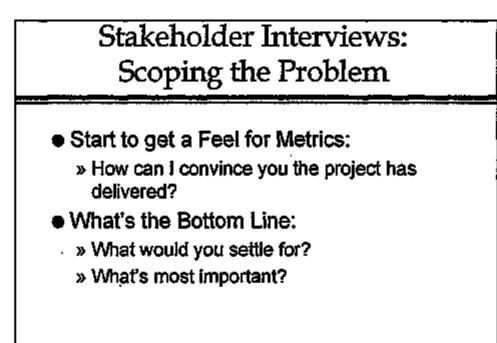
Goal: Figure out the Politics, as best you can

- Draw an organizational chart, with everyone involved in the project.
- If there are other companies involved you need to chart that organization(s) players also.
- Review and assess the Players:
 - » Who is powerful, listened to, gets their way?
 - » Who could ruin your career?
 - » Who has money and interest?
 - » Who wants the Project to Succeed?
 - » Who wants their pet technology used...



Stakeholder Interviews: Open-Ended Questions

- Listen, Take Notes, Don't Argue or Sell: Listen, and ask questions just for more information, clarification
- Encourage Daydreaming:
 - » What would a perfect solution look like?
 - » What is really desired? (not how, what)
 - » How would this make things better?
 - If appropriate, show similar systems, demos, etc.



Creating Order out of Chaos: Matrix of Customer Desires

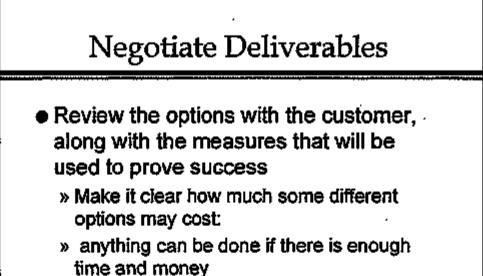
- Brainstorm with customer group if available
- If there are customers with very different needs, create a list of desires for each customer from your interview notes
- Create a Customer's desires matrix, noting who cares most about what

Document Customer Desires as Measurable Objectives

- Example: Instead of "User Friendly":
 - » "Novice can use the system to do x after 30 minutes of training"
 - » "Users with more than 1 hour's experience make less than 1 error per12 major operations as described in the Acceptance Test

Find Common Priorities & Plan Strategy

- Review Complete Customer Desires Matrix with all the Customers: Find Overlaps
 - » Ask for rank order requests
- Quality Functional Deployment (QFD): How will we deliver?
 - » What's technically possible: what will it cost?
 - » Where's the biggest payoff/risk?
 - » Create cost/options estimates for approaches
 - » Determine our presentation/proposal plan



Create the Requirements Document

Write the Acceptance Test Before Development Starts

Write the Requirements Document

- Write out the Acceptance Test criteria for each requirement:
 - » this defines exactly how the requirements will be measured
- Review & Renegotiate the Requirements and the AcceptanceTest Doc. with the Customer
- Create a Detailed System Test in the general design phase as implementation details arise

The User Interface is Defined/Refined during Proposal

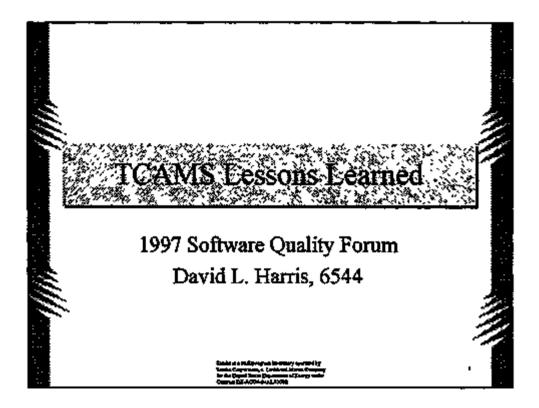
- Prototype and review the initial user interface quickly (Reusable code?)
- Use the people who will actually be using the system for the user testing:
 - » They become champions for it's acceptance.
 - » They know their jobs, and how it will be used
- Complete the User Manual before coding the User Interface: It's the Requirements Document & Acceptance Test for the UI

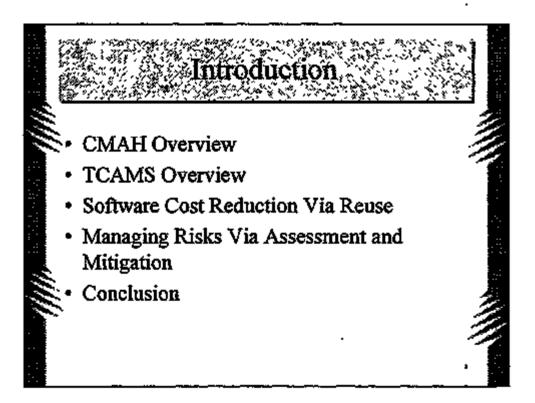
Rapidly Prototype the whole system if possible

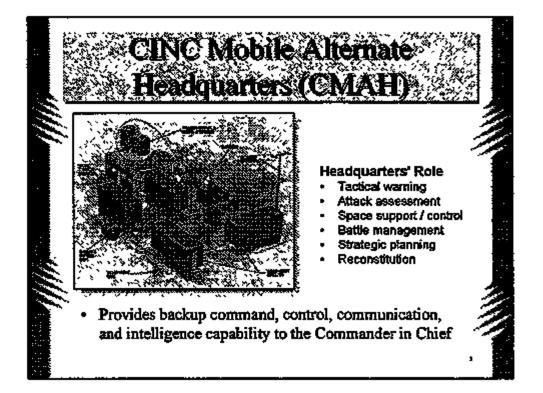
- Reusing a configurable system increases robustness and cuts development time
- Demonstrate and Modify System as Requirements are renegotiated.
- If New Development, Build Incrementally
- Structure the Project with Many Milestones: coordinate incremental changes to deliver new functionality

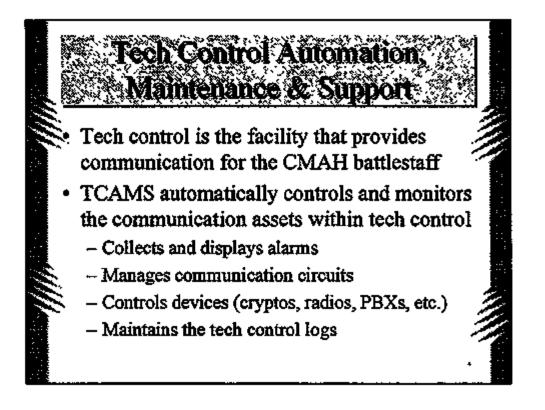
Summary: Listen, Document, Review, Update

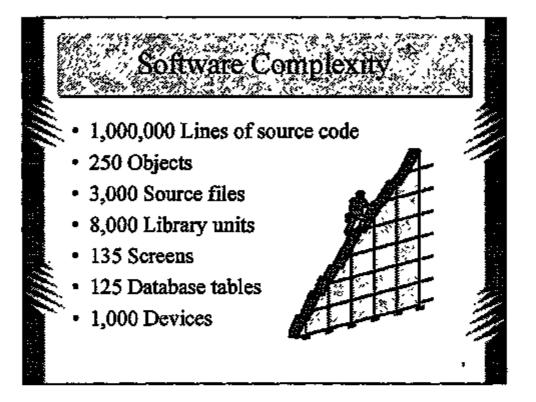
- Find out who the customers really are: Organization Chart Review
- Find out what the customers want: Customer Interviews
- Figure out what the project needs to deliver: Customer Desires
- Figure out how the project will deliver: QFD, Proj. Pian, Child Design Matrix
- Document how we'll know we delivered: Acceptance Test, User Interface Manual & Milestone Reviews as the Project is Implemented

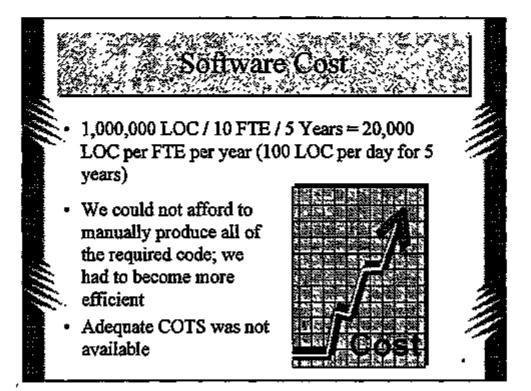


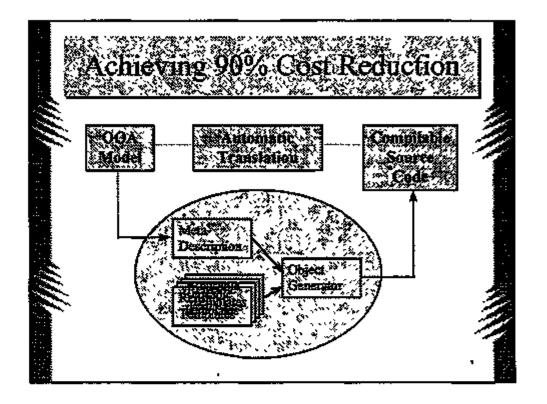


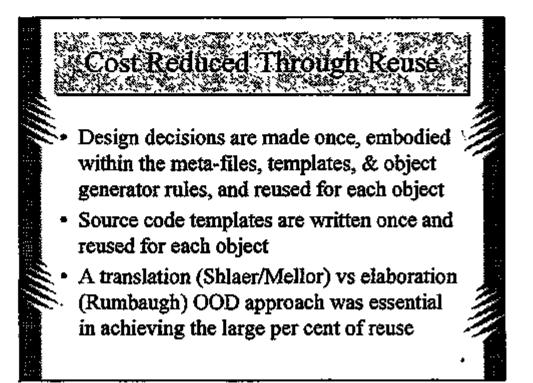


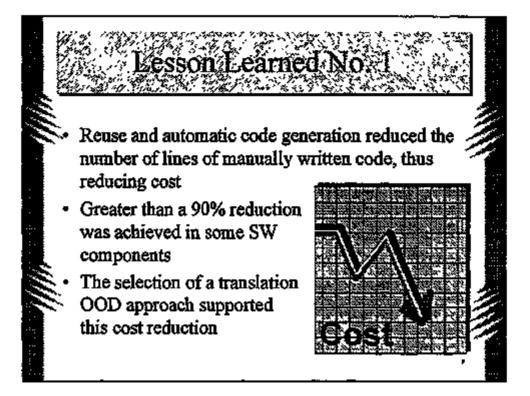


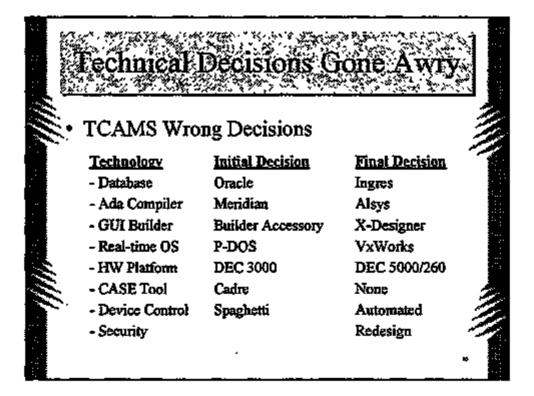


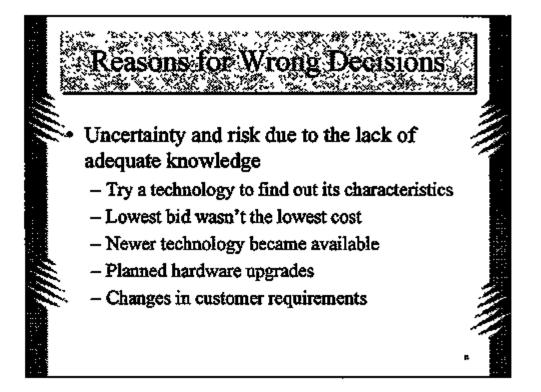


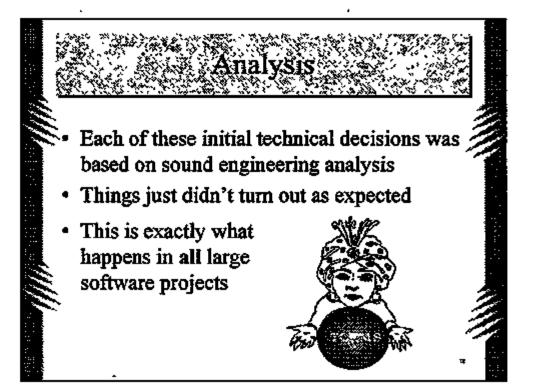


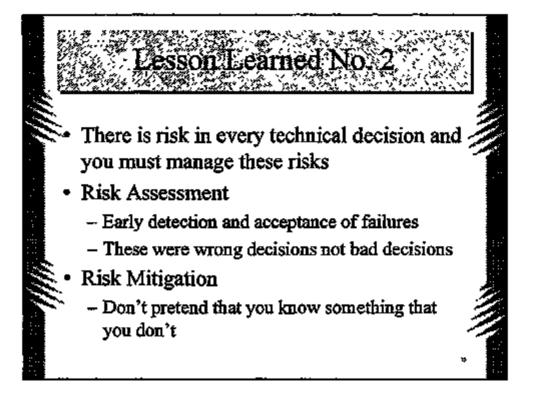


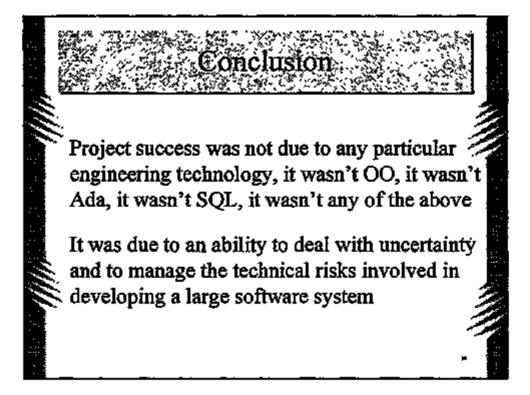


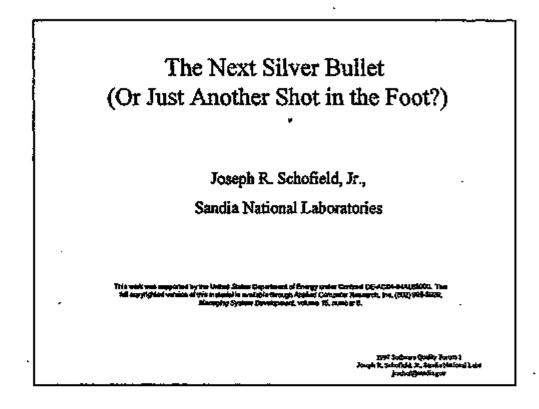


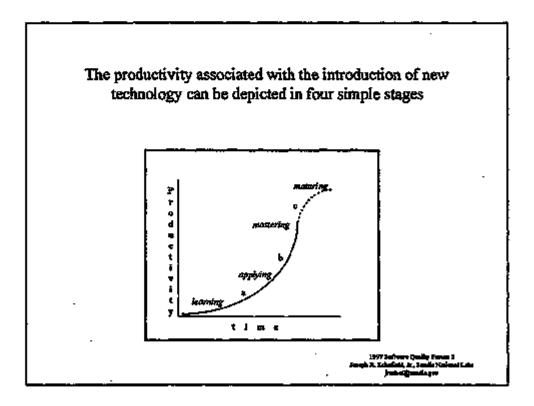




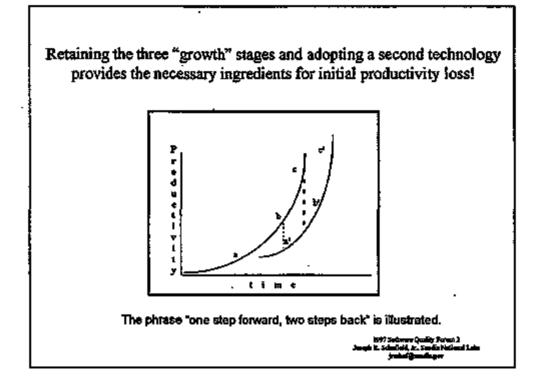


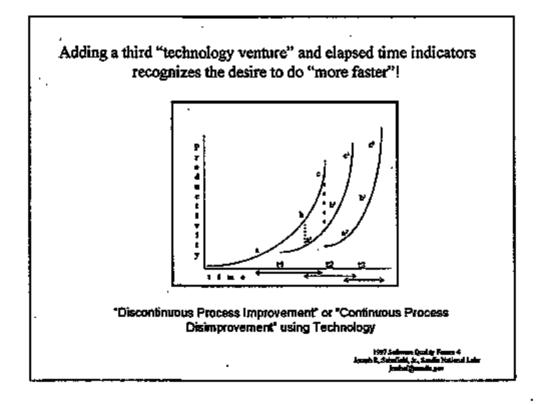


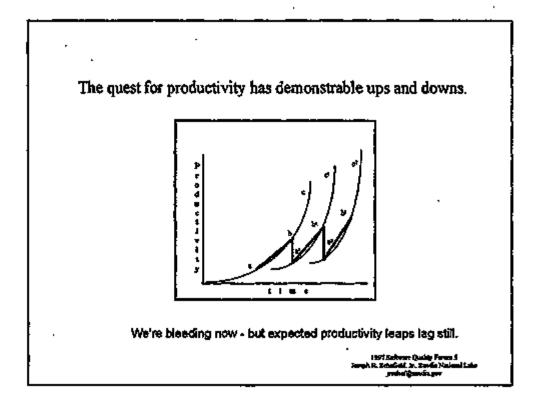


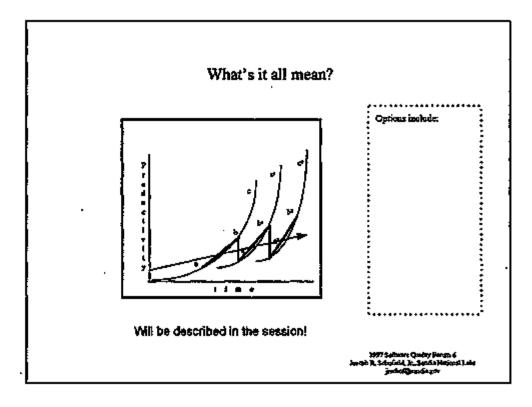


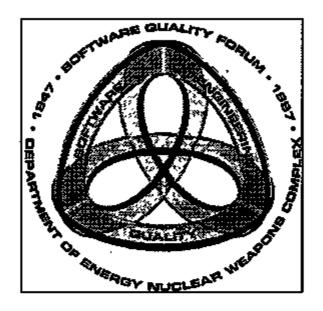
I











Session B1: Software Testing

Chair Larry Rodin Pantex Plant

.

Session : Paper #	Author(s)	Title
B1:1	Debra Sparkman Lawrence Livermore National Laboratory	A Working Testing Process
B1:2	Nancy Storch Lawrence Livermore National Laboratory	Testing the Design and Operations of a New Badging System
B1:3	Dwayne Knirk Sandia National Laboratories	Establishing a Three-Way Agreement: Specification, Code, Test

ί,

1997 Software Quality Forum

A Working Testing Process

Debra Sparkman Lawrence Livermore National Laboratory

April 1997
This work was performed and/or die US Department of Energy by Lawrence Livermore National Laboratory
ander Contract Na. W-9405-Eng-42.

Argus Overview

rgus

1

2

t

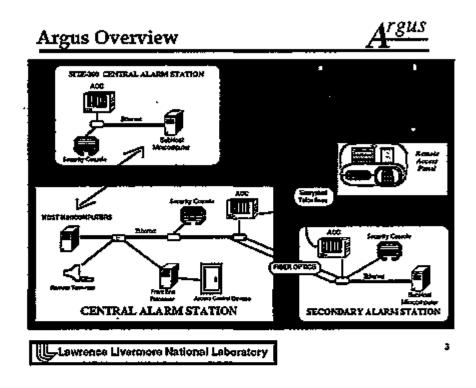
rgus

Automated Security System

Intrusion Detection

Access Control

Dispatch Center



Developer Testing	Argus

Q Unit/Package Testing

C Ada test packages for shared modules

Integration & System

C Performed on development system using "mock" utilities

Lawrence Livermore National Laboratory

__}

Conducted on separate system

- Based upon configuration of all customer sites
- Physical equipment in most cases
- Flexibility to configure system to allow parallel testing for different sites

rgus

Focus on regression testing and new major feature testing

- manual testing
- 🗆 repeatable

Lawrence Livermore National Laboratory

Independent System Testing cont. Argus

- Test planning based on priorities
- Time allows, perform special feature and defect correction testing
- Test anomalies tracked and reviewed by Test Leader
- Test summary
- Major coordination efforts, frequent meetings with development staff
- Q. Metrics collection of cumulative failure profile.

Testing Process Tools

C Test procedures priority & pass/fail log

rgus

7

Test incident report

Test sequence log

Lawrence Livermore National Laboratory

Test Procedure Priority & Pass/Fail Log

C Excel spreadsheet

Testers input data during testing

- 🖬 start & stop times 🗉
- pass or fail status
- O initiais
- 🗆 comments
- Automatically calculated fields a duration

Status reporting

testing completed

C testing to be completed

	Excel	le	et						A ^{rgus}									
					ì	1												
(B) H	fsi hadin Ch	kal (hadini	H	i iz Dian	1	:	b h	20A	6	Nijur Kante (in	utui Mair Ja	ityai ku (ki)	1	ſ	li	R,	ļi Ē	 Caurds
H			ļ	ж							6JT	認						
ZB,	lensi Si Sensi Koji (k) sesi k. Mijanići ji k Niši Senji Kništaram.	Rèner lite;	1	Ħ			1065	ikter i	L3	ið	8		1				3	
26	iztok i lesklaigistasi sinilapiin bis Map nittisesi		I	12						1	1	I				1	3	 ವಾಘನ
26	icaty i k 26 kpc nainy ink analysia de site ing	iie is	I	ti			tata.NST	22	ttill	5	ß	¥3	ł				5	

Lawrence Livermore National Laboratory

Testing Process Tools

rgus

10

Test procedures priority & pass/fail log

Test incident report

Test sequence log

Test Incident Report

C FileMaker Pro Application

C Real-time defect reporting

All defects are collected

- C Software
- 🗅 Hardware
- Q Test Process (tester error, test procedure defects)

Å

;

6

rgus

11

Test Configuration (test system specific data)

Test Incident Report cont.	<u>A^{rgus}</u>
Two Impact Categorizations	
C Testing impact	—
O Release impact	
Status & Approval Signatures	
D Assigned to	
C Resolved by	
C Relested by	
Approved by	
	· _
. ·	-
Lewrence Livermore National Laboratory	12

eellast 19	Test Insident Teaking System	•
	ATT 22A HYBRID	Test Log 124
in Same_12/18/96	THEORE APPTCHTESTER	i
wytersteinen Persotti	وحولا المتازز لملخ ومعلا	parkmen 1/10/97
Ale and a second se	Secondaria	
Gaft Water	Reptatable	
		Manager and Carp
with example law will be track on the Trive	ann († 2522-7-1533) (Pranton i na de Parlieux Inda, fit dans parts (Pranton de Altr, St Maad, fat Parla part (Pranton) Maad, fat Parla part (Pranton)	
Fallgard in sell in an an it i 11 yr an de fynnyn a'r arl ar i 2anael 3 Nawr yw yn yn yn yr ar yn ar yn ar yn	helds, fy dewy przew sprewskaw (ż. + ATP. 3) własz, świ Piała w 28 ard 37 Rał nate Zryne	i Menda Zin Balera Jajena
Failgand in seil i seil i seil an ib à 11 yi antari 1970 y si d'art la mart 2 Biann 15 Tari lag Bigar (1 - Milan Tari lag Bigar (1 - Milan Tari net malana vi da seil-dast to	helds, fy dewy przew sprewskaw (ż. + ATP. 3) własz, świ Piała w 28 ard 37 Rał nate Zryne	
Singer in sett i ser an in i si si anter i yr ny sett i ser i ser i se i sett gen yr ny sett i ser i serner i anter i gen yr i Singer i In set yr i ferni sett i sett i sett i sett i <u>an sett yr i ferni sett i sett i gen yr i sett i sett i sett i sett i sett i sett i sett i s</u>	helds, fi dawy pages proverbans (2 = AT2, 2) stand, for PAN to 28 Ard 37 Fall safe Zeryne Fling writi siwel. I'm ATP f. Th. ManZ, to be principal day instructes traded () years bell, safe of the s Top cash the by ATP 3 (2) we day.	a tim um tilguas tim 17 a sez set and 1. Norda 7 ja Bulsen Jajann
Management has well a sear on the origin and and typing well a sear 2 annual 3 for the forguna of the search of the for the forguna of the search of the forest the search of the search of the forest the search of the search of the forest the search of the search of the search of the	initia () down param terretakan ta' ATT, 35 Anna, fer Ann an 28 and 37 Ral nata Zeyna Along weld sinet. (in ATT f. Sa. Manda ta ha ye seriartake itu (in ATT f. Sa. Manda ta ha ye seriartake itu (in ATT f. Sa. Manda ta ha ye seriartake itu (in ATT f. Sa. Manda ta ha ye seriartake itu (in ATT f. Sa. Manda ta ha ye seriartake itu (in ATT f. Sa. Manda ta ha ye seriartake itu (in ATT f. Sa. Manda ta ha ye seriartake itu (in ATT f. Sa. Manda ta ha ye seriartake itu (in ATT f. Sa. Manda ta ha ye seriartake itu ta	a tim um tigan tim () a serest and 1. Norda Zin Jadem Jajenn 1. Norda Zin Jadem Jajenn 1. Para antida antigar ()
Algorith for sell - net ar it is 11 y and and by regardle and the sell of the sell of the sell of the sell in the forget of the sell of the sell in the sell of the sell of the sell of the sell in the sell of the sell of the sell of the sell in the sell of the sell of the sell of the sell of the sell in the sell of th	ndde († deur pages pressions das ATT, 35 stend, for GAB to 28 and 37 Kalante Segur eijeg until sinet. (in ATP f. The Mark to be printede dur in ATP f. The Mark to be printede dur (in ATP f. The Mark to be printede	a tim um tigan tim () a servet and 1. Norda Zin Balera Jajana 1. Norda Jajana 1. Nord
Millioned September 2011 - Series are the a 11 yd handberd Syrring weld dark Lawrened S Series Gregories - Schwa San met sweldene wilde ar Statistic fan handeldene (1991 / Frankling were sweldene ar AS San Statistic - Statistic ar AS Plant Frankling Analysis tar Ed	ndde fi daw sam i sreetiken it - Aff. 35 stand for Ran to 28 and 37 Falance Zeyne Sin AFF f. In Mark to he prioriale in the state in the state in AFF f. In Mark to he priorial it is in AFF f. In Mark to he priorial it is in AFF f. In Mark to he priorial it is in the state in the state in the state in the state in the state is in the state is in the state is in the state is in the state is in the state is in the state is in the state is in the state is in the state is in the state is in the state is in the state is in the state is in the state is in the state is	t Marda (m Salem Ja) nov Marda (m Salem Ja) nov MTP FIP menytakim, Tenymiji) 12/39/96

Lawrence Uvermore National Laboratory

Testing Process Tools

rgus

13

14

- Test procedures priority & pass/fail log
- Test incident report
- Test sequence log

Test Sequence Log

Provides a sequence of activities

Giobal viewpoint of product testing

Logs software products version & date identifiers

rgus

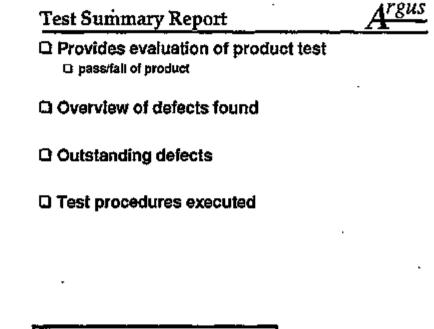
Lawrence Livermore National Laboratory

Test Sequence Log

Pale and Time	Test Procedure 18	Converts
041210, 1548		Recent or your stight for twilling
961213, 1511F	Youl Svalors	1. APP 2.34, 4- Day, 16:17:02, abackum 25of (APP_650) (Master APP)
•	20 terms	2. APP 2.16, 30- Co. 22.43, shedrown Och9 (outbears, slove) (subsers 4APP)
	Cordian atom	2. AP2:24, 4 Des 95, 15:17:45, sheet over advect (hybrid)
		4. Argue VAX(Tools 6.4 (4 Dec. 12:34)
		5, Argus Tools 8.4, 10-max 1896 10:00
		5. BXLA3.27 3 Hout 4.5, 15- Nov, 11.31
		7. Rt annuty 1.1, 27- Aug. 10:58
		5. GAN 99.57 detect 12- day 1995, 17:19.
		9. Censel e 2.Ch, 11- Nev, 16/29
		10. MPC3.04, 4- Oro-1996, 11:36
		1 * 61.45 How 19.7. 21+ Nov. 00:29
		12. O4U 2.9, 11+ 3m, 11:34
		13. CCTV 8w ver 1.7. 21- Dec 1995. \$5:52
		14, Phonebook Berver 1,5, 22-Doo, 19:13
		16. Yuna Ada 1.2, 22- Dav 1995, 98:39
		<u>hə.vəru 1 4. 19-50, 1949</u>
61213, 15#5	Tes bes/ server	Four bour earsest racia waring: [[Hybrid] annear rack], [AFP, Fahaad' Waleye test bou], [PTU
	radi wiring	Microf + Text taxe) },
Mi1213, 1620		Chaptic get gyret gen Sunight grant i ty .
61213.16:25		Europedid test incorrect mice for the day
M1214_0740_		Pages (next program at less the starting.
061215.08-00		Rear telling
61215.12:10		
61214, 14:20		Pages Maid Spail and
961236, 15:15		Next is nervew name is all tasks with Foot and Brock before a dust mining it several failance
,		Piotes fasta pages falles
P61214, US29		Suspension setup for the day.

Lawrence Livermore National Laboratory

16



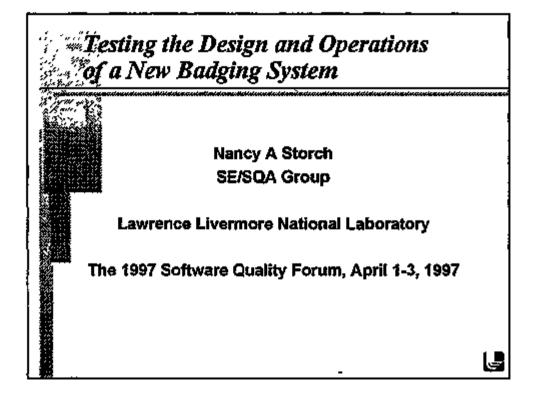
Lawrence Livermore National Laboratory

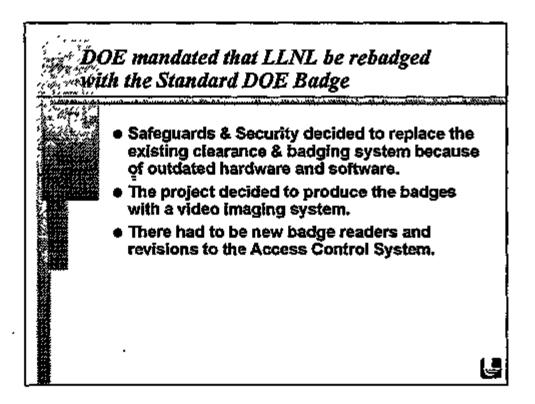
ų

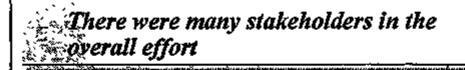
Summary _______Argus O Manual process Repeatable O Process tools O Defect tracking Defect tracking

Sevence Livermore National Laboratory

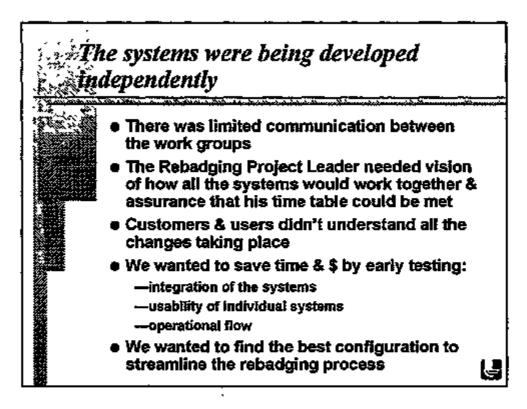
17





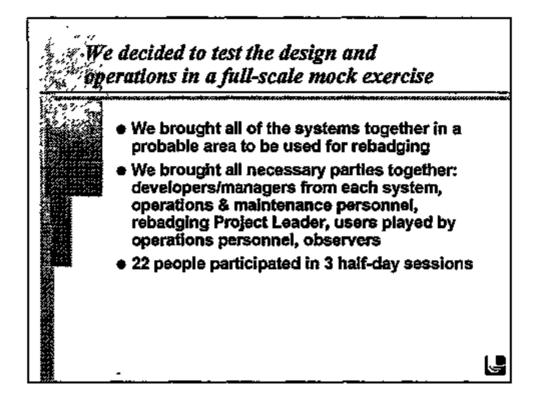


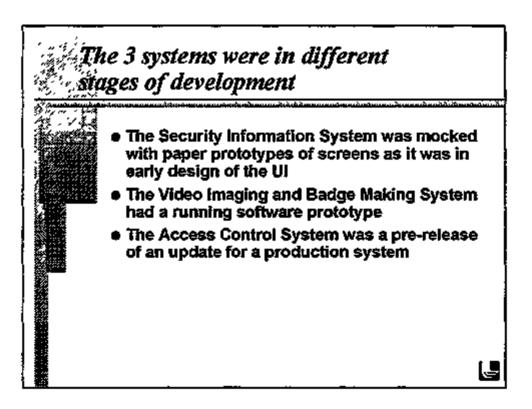
- Badge Office (BO) & operations
- Central Clearance (CC) operations
- Security Information System developers supporting BO & CC (reengineered system)
- Video Imaging and Badge Making System (new system)
- Access Control System (new release)
- Rebadging Project Leader

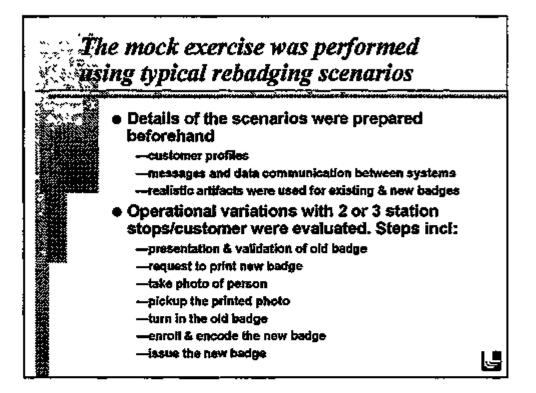


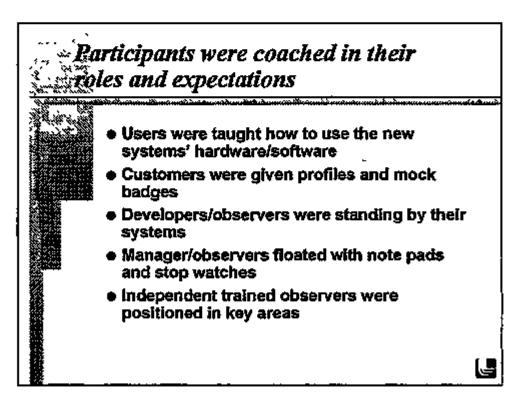
¢'/

G







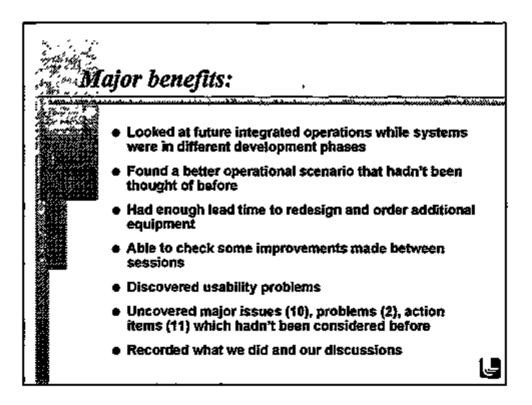


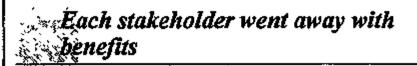
Learning from each session was applied the next day

- Each session started in the lounge area with an explanation of the scenarios we would be testing. Roles were assigned.
- After scenarios, we gathered again and collected observations, recorded metrics and did some analysis.
- A facilitator compiled lists of issues, problems, and action items which were added to with each session.
- A plan was made for the next day based on what had happened. The day's activities, questions & comments were recorded.

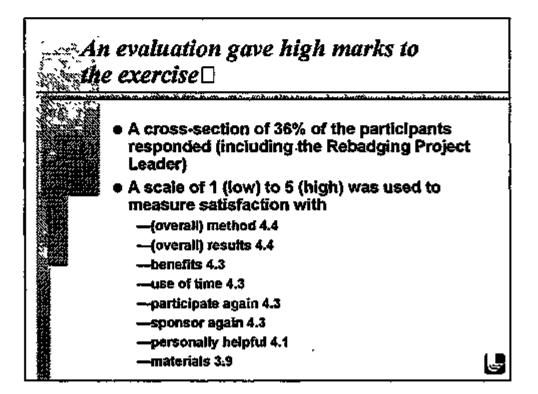
5

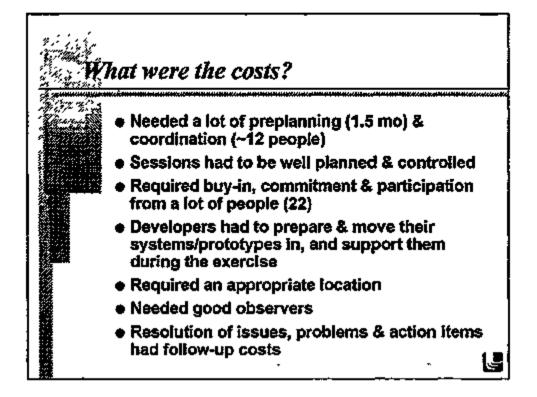
We held a final concluding session



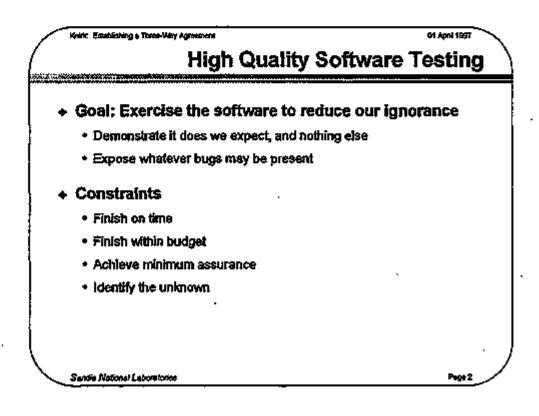


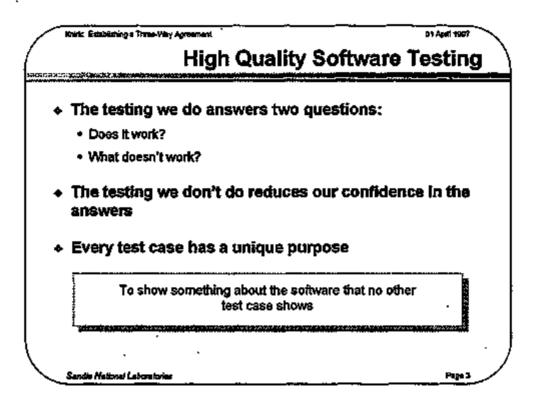
- Badge Office decided to look at another badging location & do more mock exercise s with other badging scenarios, resulting on operational changes and remodeling
- Central Clearance (played customers) became familiar with their sister organization, the Badge Office
- System developers uncovered misunderstandings, erroneous assumptions, and omissions
- Rebadging Project Leader learned more about the systems and operations, and gained confidence that rebadging would work

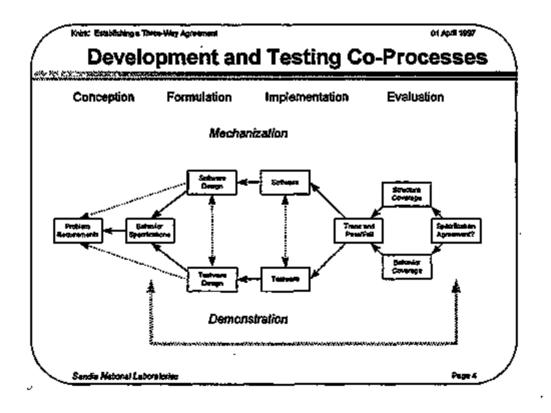


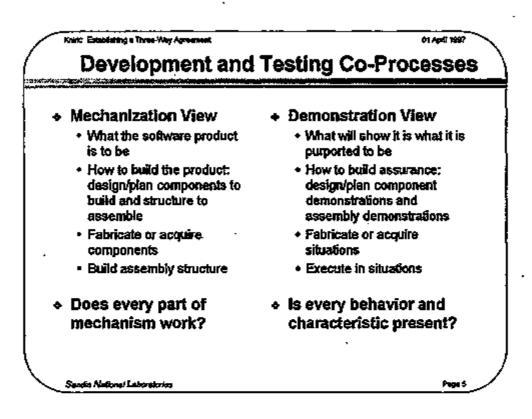


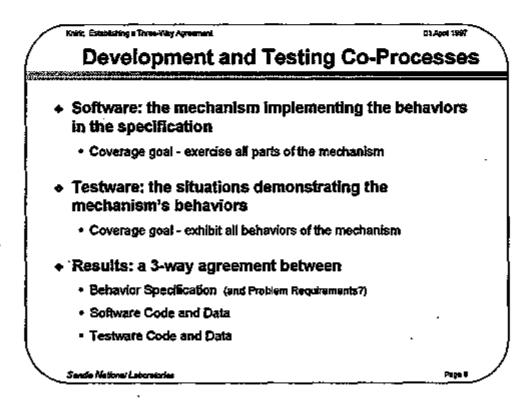
Robit: Establishing a Three-Wey Agreement	O'L April 1997
Establishing a Three-Way Agreeme Specification, Code, Test	ent:
Software Quality Forum Albuquerque, NM 1 April 1997	
Presented by Dr. Dwayne L. Knirk Quality Engineering Department Sandia National Laboratories, Albuquerque, NM	
SAND97-XXXC This work was supported by the United States Department of Energ under Contract DE-AC04-94ALB5000. Sendia is a multiprogram laboratory opwated by Sandia Corposition, a Locide Company, for the United States Department of Energy	
Sandia National Laboratories	Page1

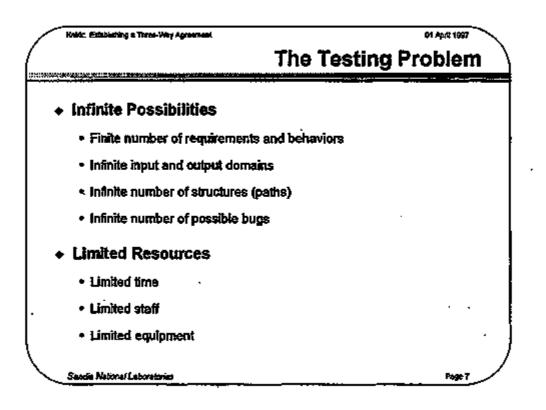


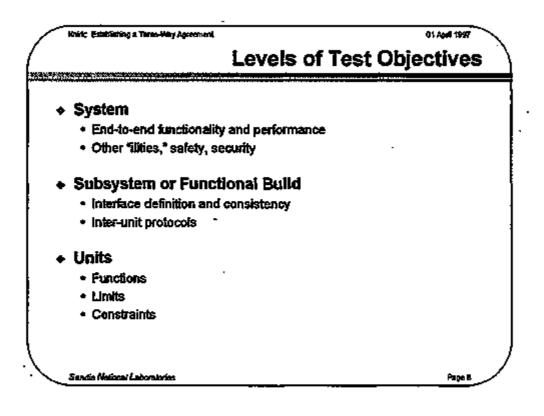




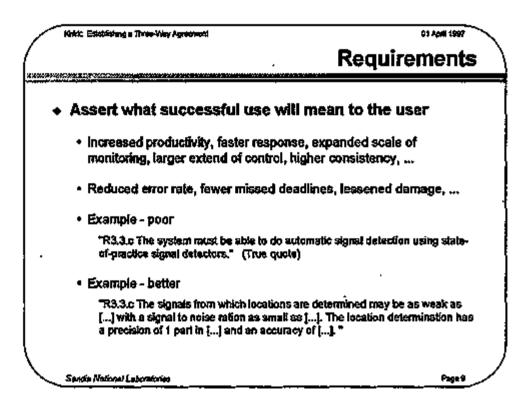


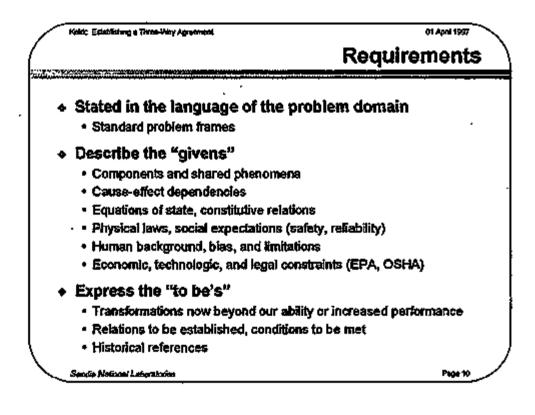


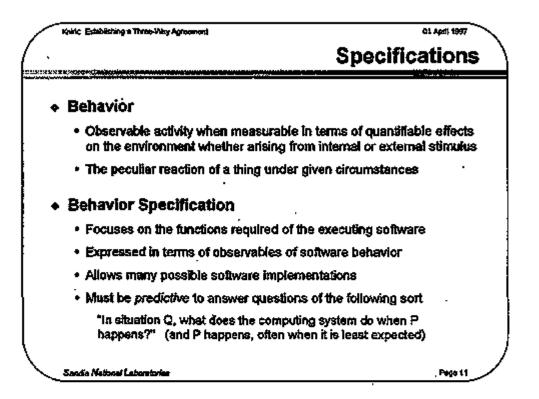


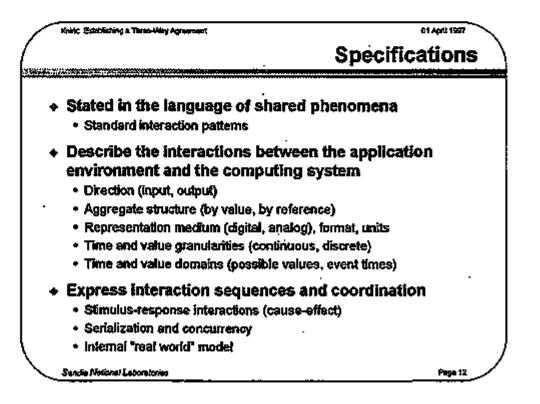


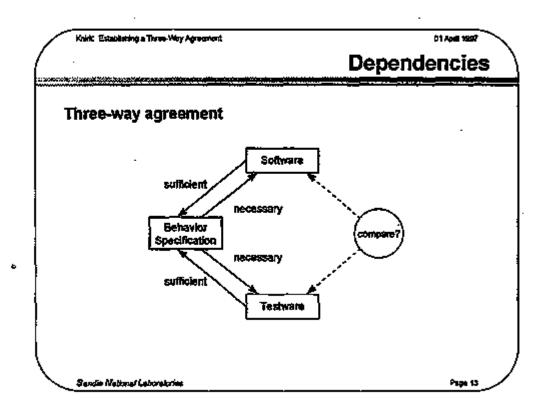
ر

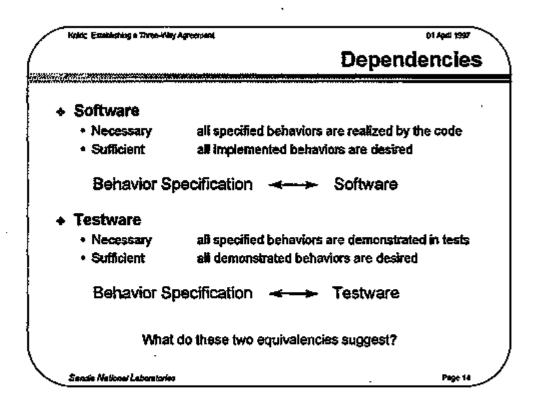


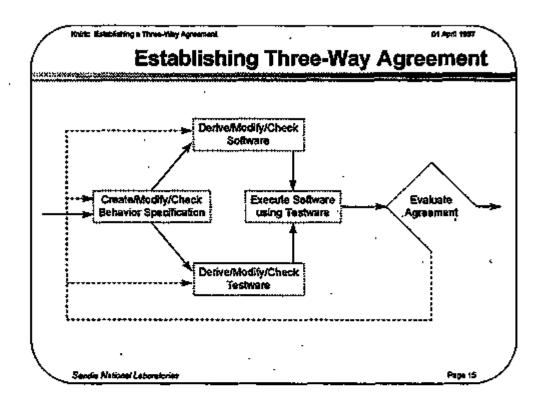


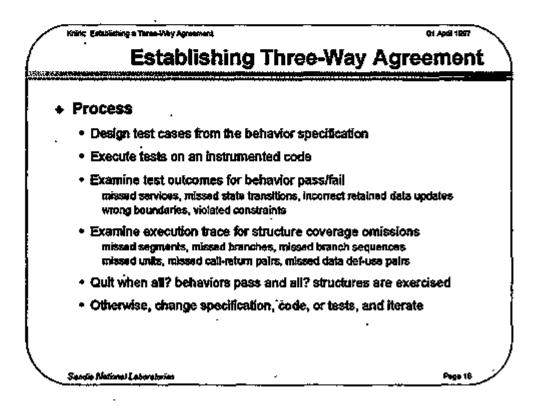


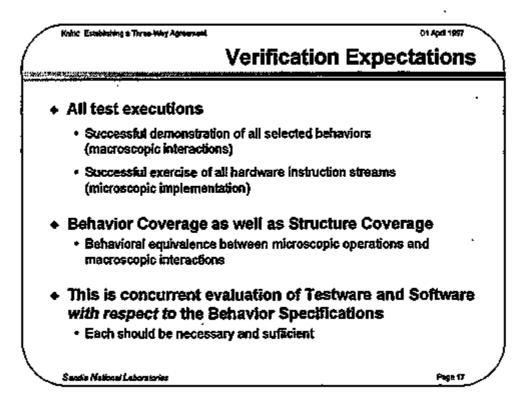


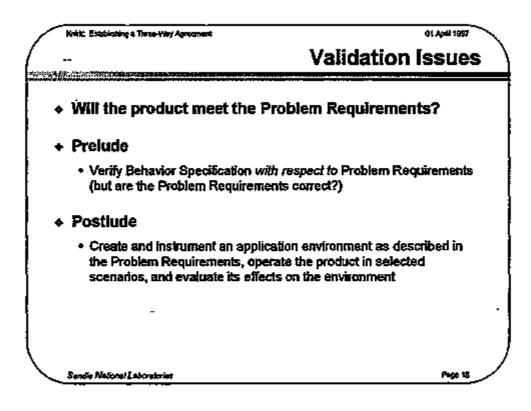


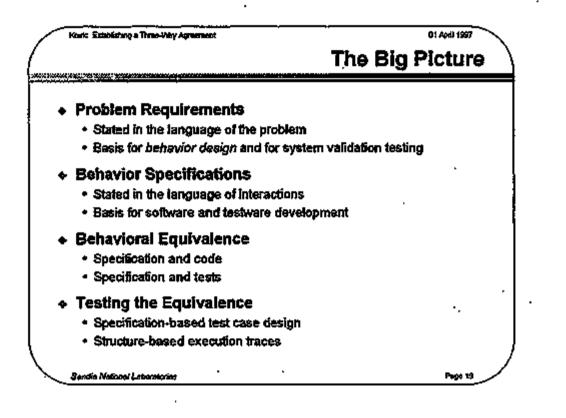


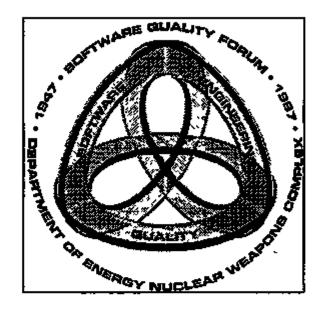












Session C1: Software Quality for Scientific Applications

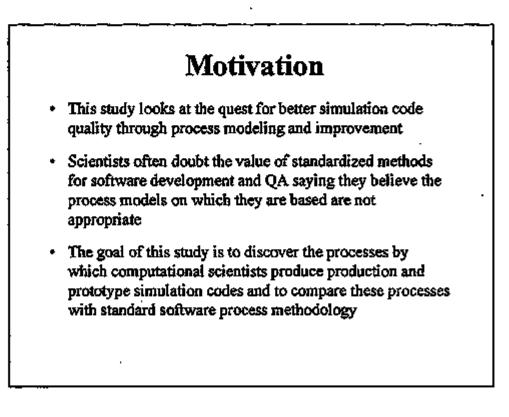
Chair John Cerutti

Los Alamos National Laboratory

Session : Paper #	Author(s)	Title
C1:1	John Ambrosiano & Robert Webster Los Alamos National Laboratory	Software Quality and Process Improvement in Scientific Simulation Codes
C1:2	Ed Russell Lawrence Livermore National Laboratory	The SQA of Finite Element Method (FEM) Codes used for Analyses of Pit Storage/Transport Packages
C1:3	Orval Hart Los Alamos National Laboratory	Software Quality Assurance at the Weapons Engineering Tritium Facility

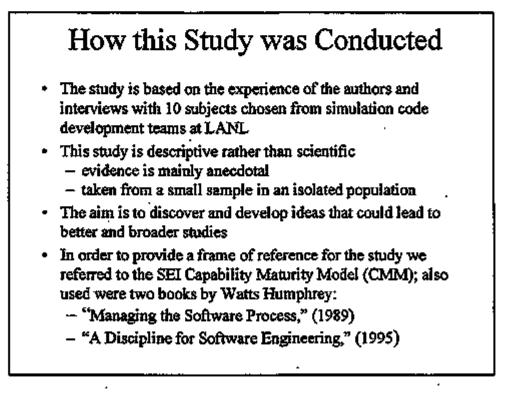
Software Quality and Process Improvement in Scientific Simulation Codes

John Ambrosiano and Robert Webster Computation Methods Group Applied Theoretical and Computational Physics Division Los Alamos National Laboratory



Background

- The authors of the study are computational scientists who have been involved in both large and small simulation code projects for many years
- The subjects of this study are scientists and computer scientists
 within the Applied Theoretical and Computational Physics (X)
 Division of Los Alamos National Laboratory
- X Division is responsible for developing and maintaining simulation codes used in nuclear weapon design and assessment
- One of the goals of this study is to try to understand our own code development processes at LANL better



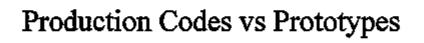
2

The Capability Maturity Model

- The CMM suggests incremental process improvement guidelines:
 - Level 2, repeatable: institute certain key practices on a per project basis
 - Level, defined: move toward uniform organizationwide implementation of practices
 - Level 4, managed: instrument key practices with appropriate measures
 - Level 5, optimizing: use measures to optimize the process

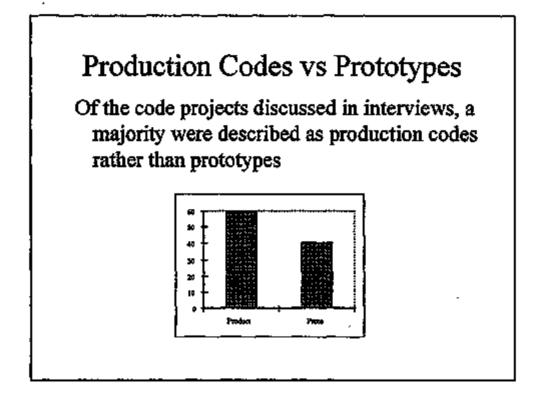
CMM (continued) The CMM suggests key practices at each level specific to software engineering Key practices considered essential to reach level 2 are: Requirements management Project planning Project tracking Subcontract management Quality Assurance Configuration management

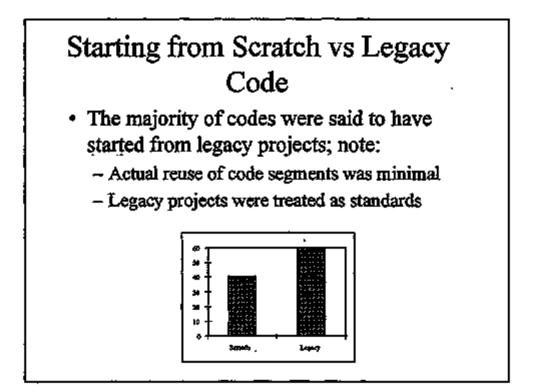
General Statisitics Project size: between 2 and 15; average 6 Many projects described as ongoing for years (1 to 15); average 5.5 Numerical application domains covered: hydrodynamics, radiation transport, neutronics, computational geometry, data analysis, electromagnetics, and plasmas Estimated lines of code: 30,000 to one million; average about 250,000



Production codes were distinguished as follows:

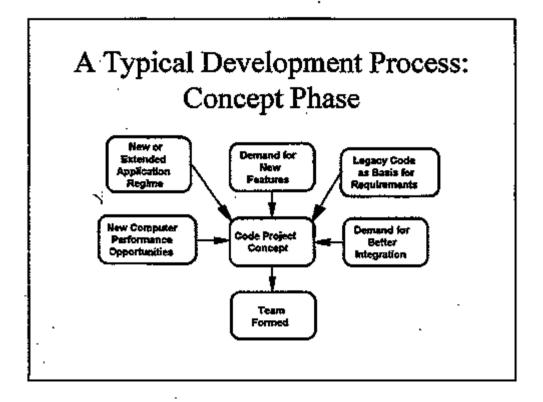
- Designed to be used by someone other than a developer
- Well documented; reasonable learning curve
- Serve as repository for models and algorithms proven to be useful; a historical archive of community experience in the intended application domain
- Give correct or expected answers to an agreed set of posed problems of practical interest

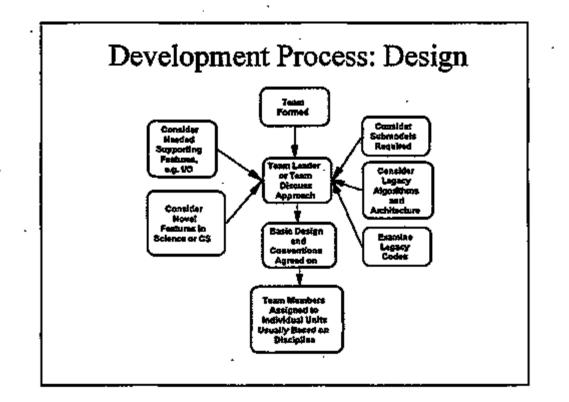




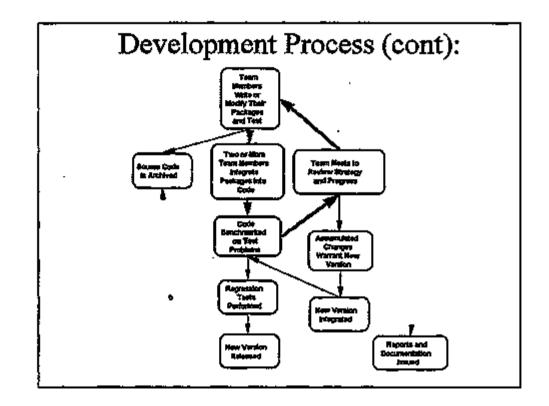
\$

.....



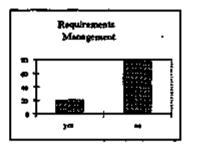


б

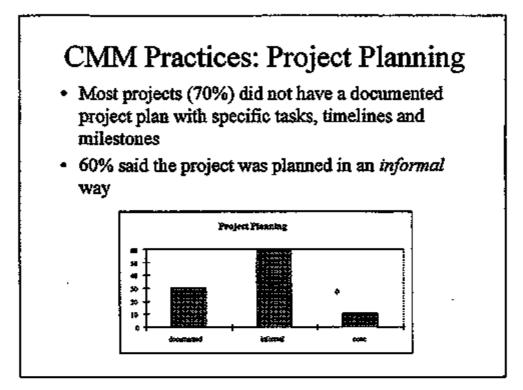


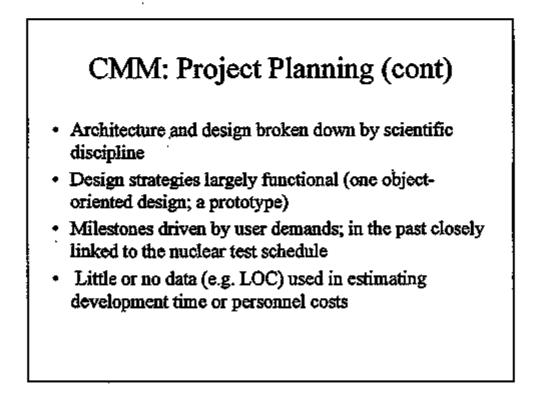
Comparison with CMM Practices: Requirements Management

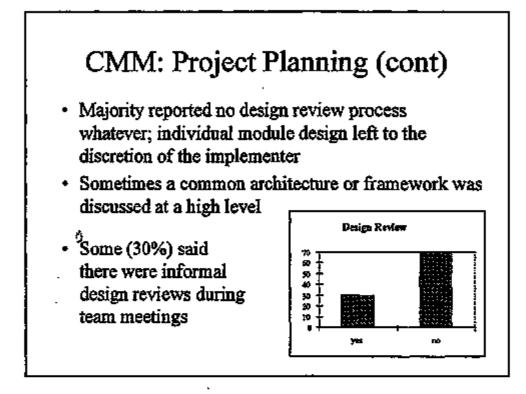
- Majority (80%) reported requirements were not developed in detail
- Interviewees told us the principle requirement is that the model produce the "correct" answer

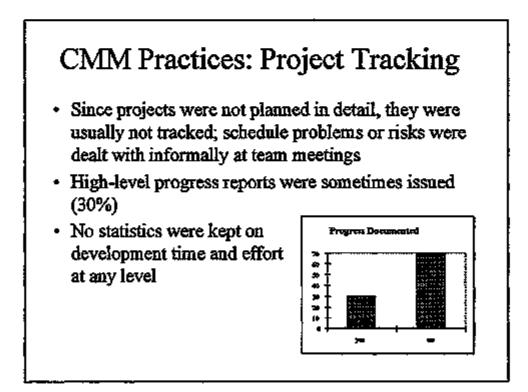


 When questioned further told us that the codes had to reproduce the results of the legacy production codes



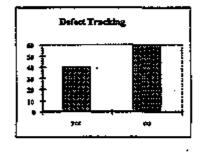




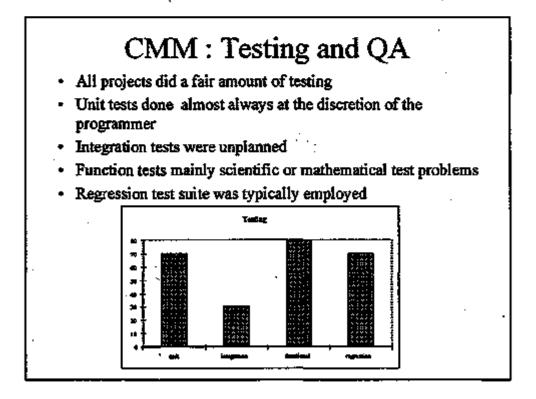


CMM Practices: Testing and Quality Assurance

• Few projects kept any statistics on defects. Those reporting some defect tracking (40%) maintained a bug report list. No statistics were kept on number of defects, type or effort expended in repair.



Ľ

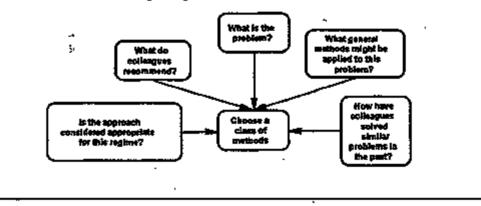


A Tale of Two Projects (cont)

- (The novel project)
 - Met its benchmarks of hypothetical test problems
 - Did not meet expectations when used on the intended design application
 - Involved new algorithms, new architectures, and new programming methods (sometimes together)
- The other project
 - Had no new methods
 - Involved 6 people for 3 years, and produced 300,000 lines of code
 - ~ Is considered a success

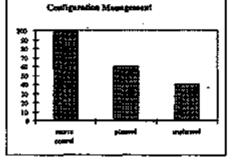
Exploration Explored: Looking at the Solution Landscape

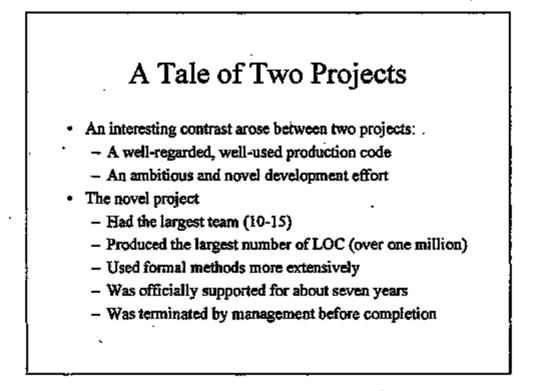
- One interview concerned only algorithm development as opposed to code development
- · The following is a process model based on that interview



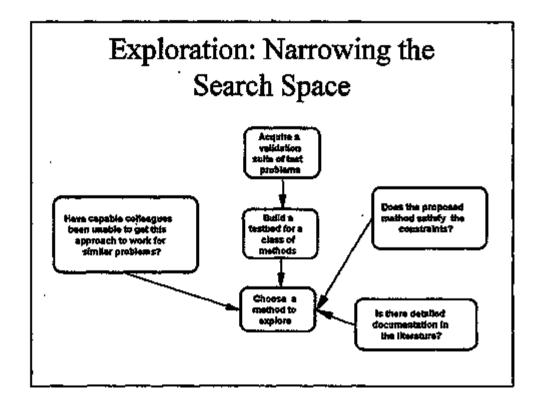
CMM Practices: Configuration Mangement

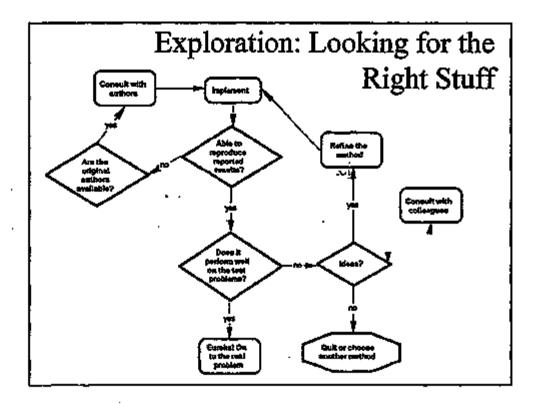
- All reported some configuration management practice
- Disciplined, planned configuration management was
 not typical
- Some sort of source control was universal





п





About the Exploratory Process

- Two interesting features stand out
 - The process is a scientific process rather than an engineering process (as it should be)
 - At almost every stage, the aim is to manage the risk of exploring unknown territory
- The difference:
 - Conventional software process models are based on process definition, process control and management of resources
 - In exploration, the emphasis is on not getting lost

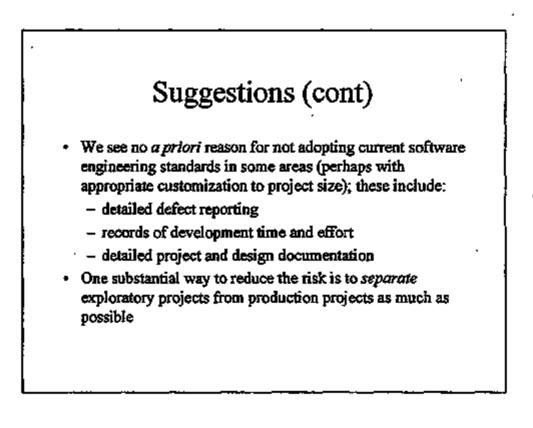
Is There Something Special About Simulation Development? Yes

- Imprecise requirements
- Higher risks in design and implementation

- The potential for open-ended testing and validation
- Strong links to legacy code
- · Relatively small project size

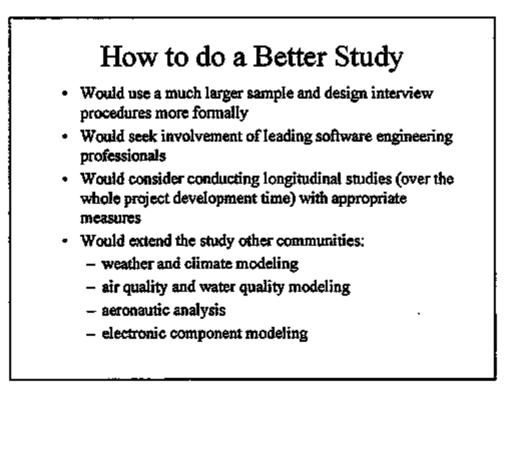
Suggestions

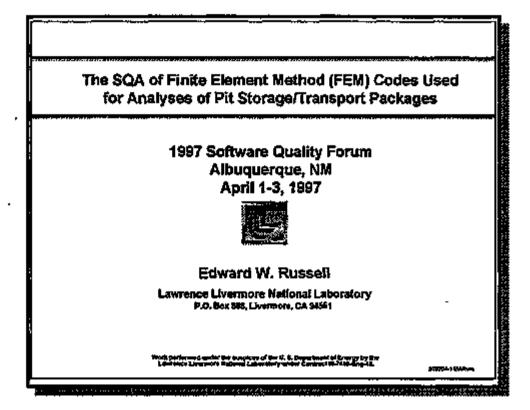
- In spite of some unique aspects, process improvement and QA guidelines such as the CMM can be of value (a substantial number of projects already incorporate CMM key practices in a weak form)
- On the scale represented in this study, one might develop some set of guidelines that falls between the CMM and the Personal Software Process (PSP) [Humphrey, 1995]
- Large organizations with several projects of this scale may be able to coordinate some generic activities like configuration management and defect reporting to advantage

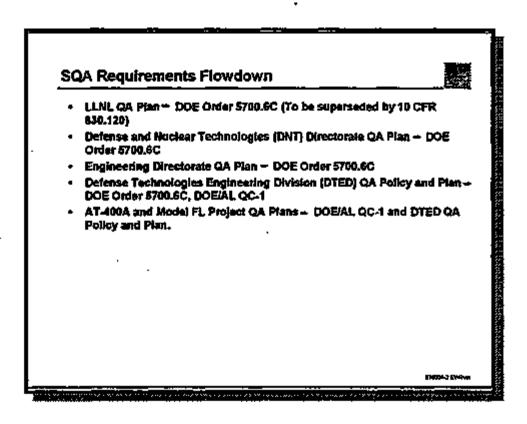


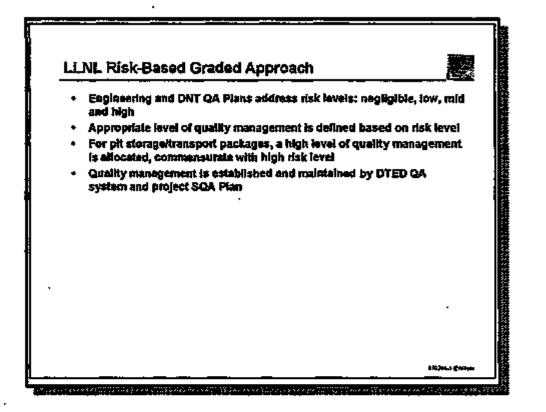
Suggestions (cont)

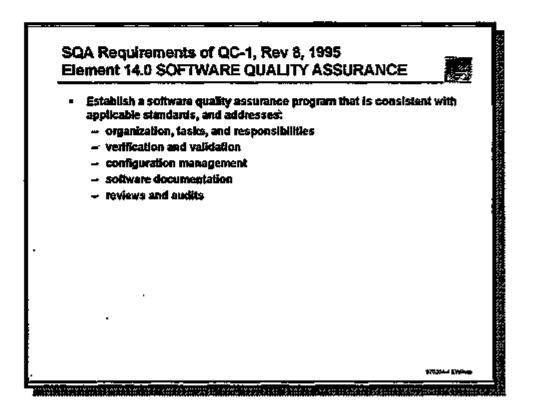
- We suggest project leaders try to nail down requirements as much as possible
- Requirements should state as clearly as possible the limits of applicability for the product; domain applicability should be defined in part by a benchmark suite of tests; benchmarks must be representative of real problems and not merely hypothetical
- Once there is a way to define initial requirements and manage changes to them, other practices such as detailed project planning and project tracking should be much easer to institute

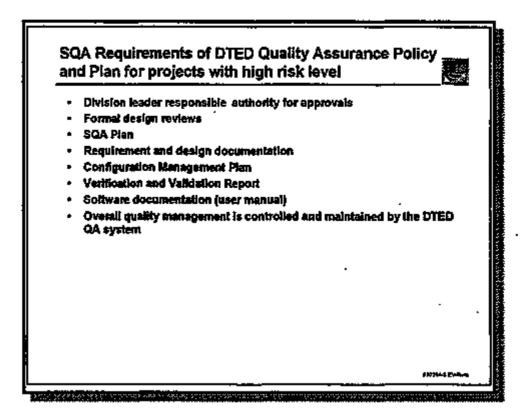


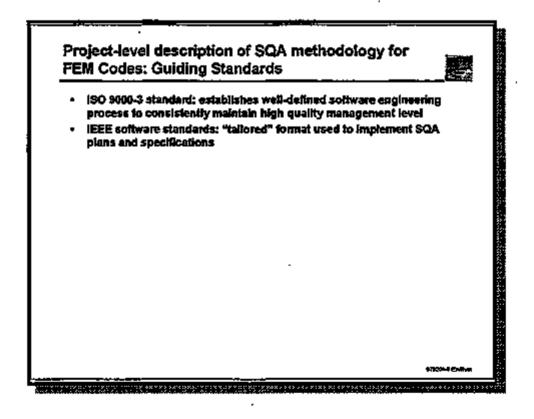


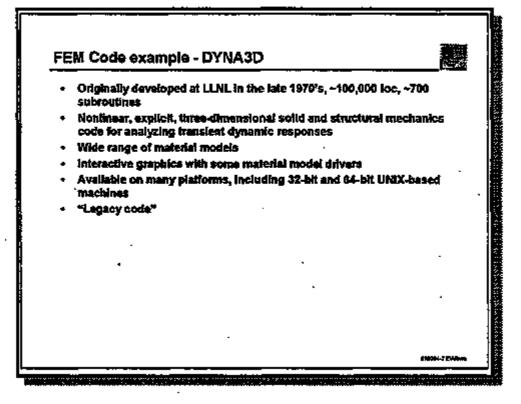


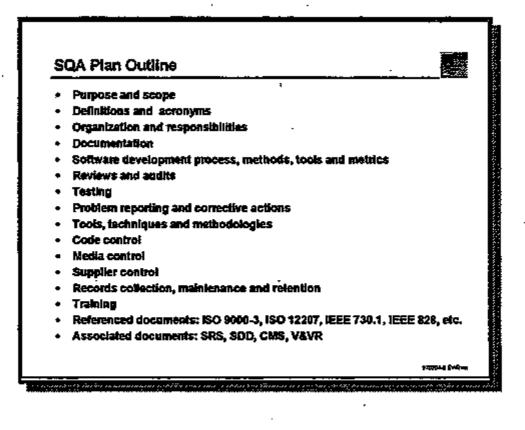


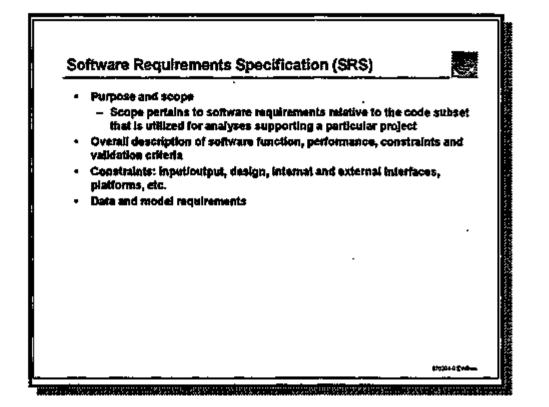


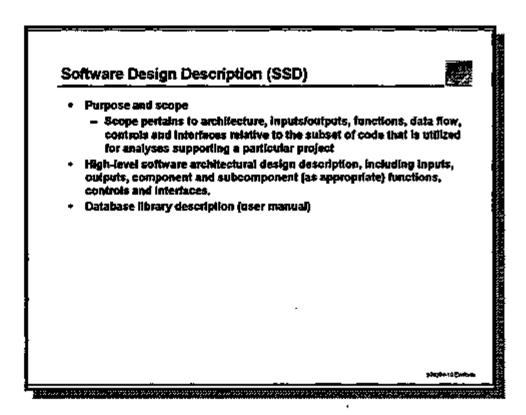


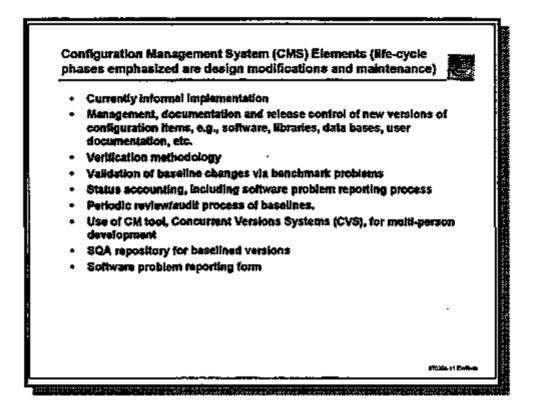


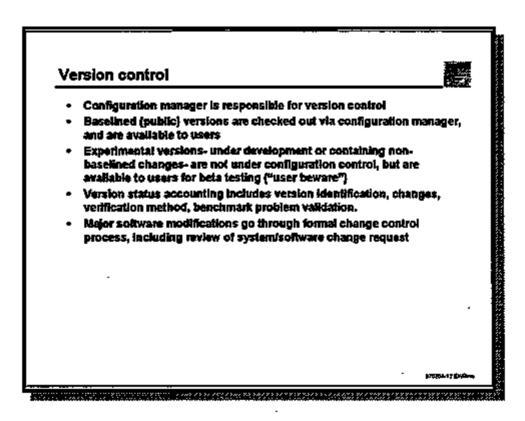


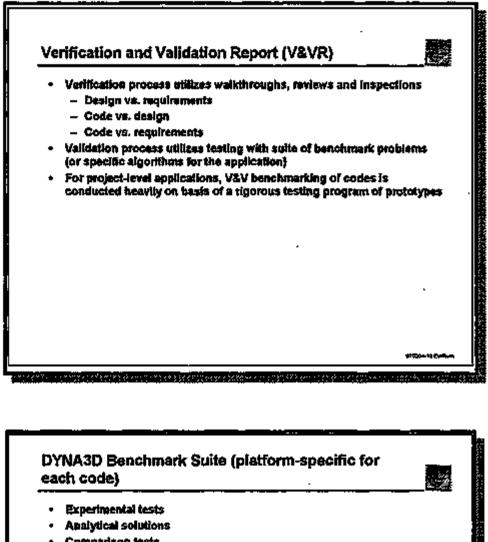


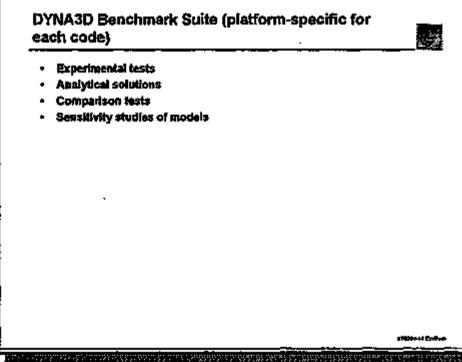


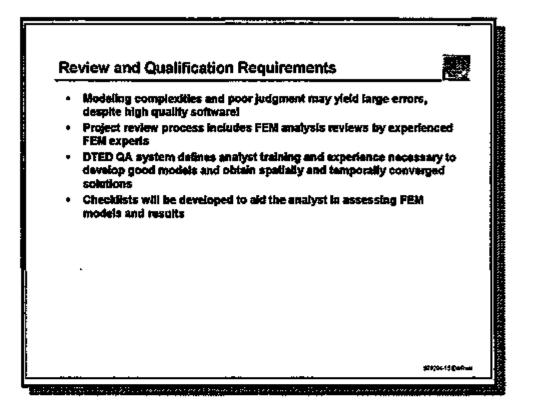


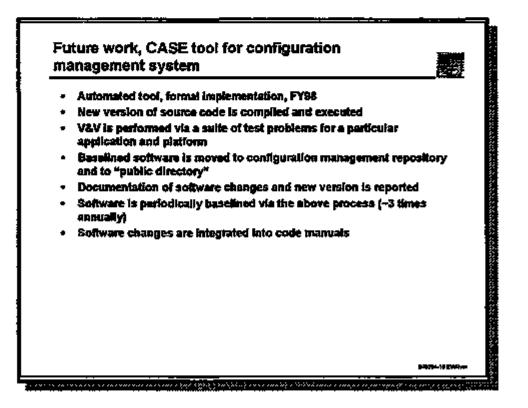






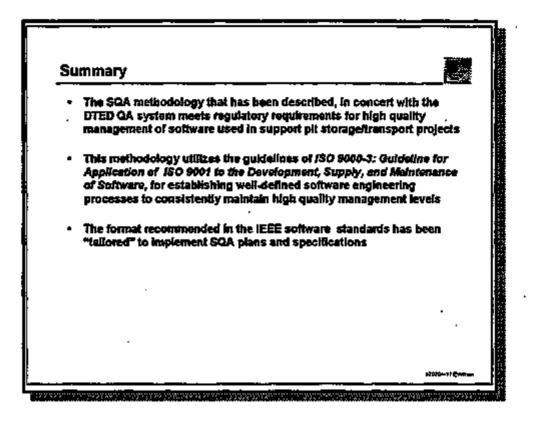


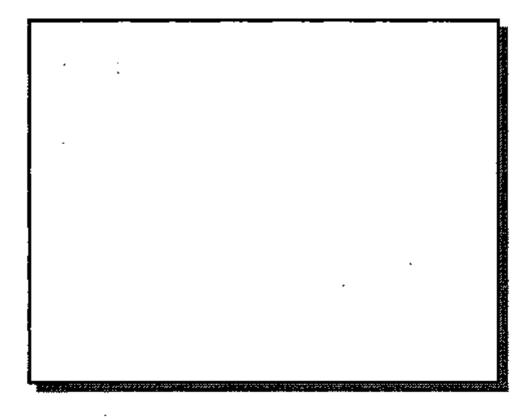


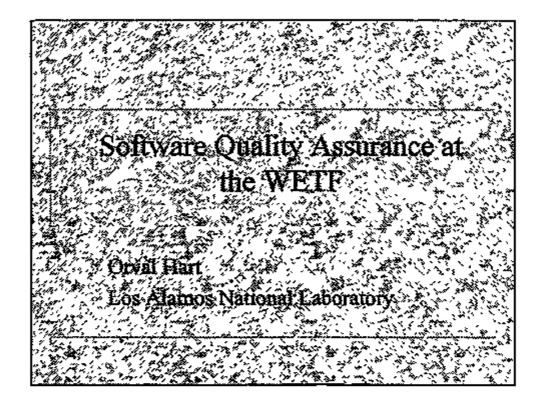


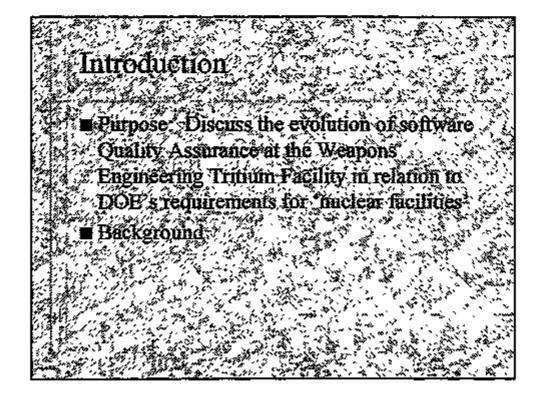
Ł

ι

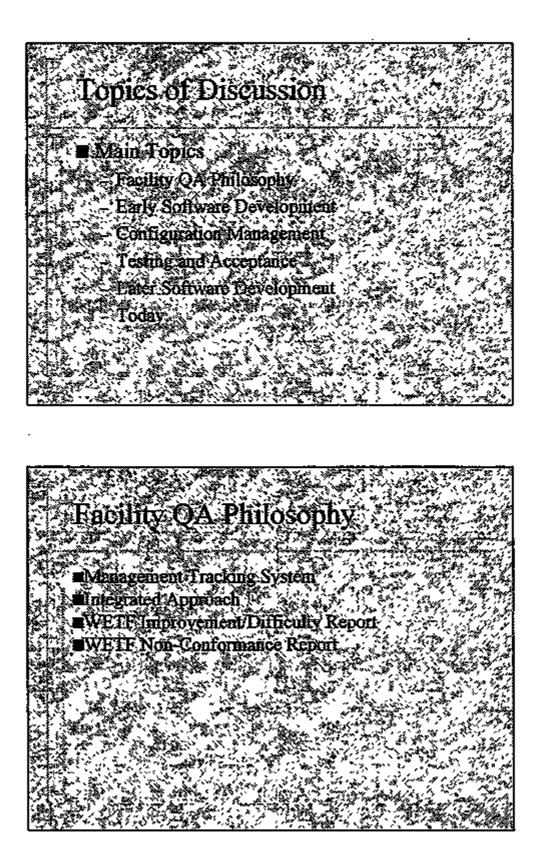








Į



WETF IMPROVEMENT/DIFFICULTY REPORT 000448Improvement/Difficulty Description (circle one) WIDR # An intermittent stack tritium months flow alarm will fail the I cs stack months stand. This, in turn, will record only "fail" on the daily stack in hysted T2 report even though the the has returned to Safety-Related? norma/, NO Signature stach Moniter Intermittent やりへ WIDR Index Description Typed Name oh) flow fails SM. daily report **Date of Report** Disposition of Improvement/Difficulty Jottem **Configuration Control?** Disposition Steps es 1 **Corrective Maintenance?** gnoch laus. Dı Non-Conformance Report ntor she SWP data Oatel RWP ile. 洇 Work Order Priority TSM Signature, Date(s) ני 2 Associated Decument Changes Functional Titlea 1. Designer/Originator 2. Section Leader 3. Tritikm Systems Qalabaata Software Procedures Drawings . Menager Building Menager ENG Division Detabase Custodian ICS Hardware Specielist Equipment Custodien 9. HSE Division Representative 10. WETF Electronic Technician. Nuclear Meterial Custodian Custodian 12. Peclify Coordinator 13. Database Designer 14. ICS Software Specialist WIDR Review/Assignment Documentation Action Required Action Checked THE Review 3 Review 1 Review 2 Seview 4 Action Completed Bγ 15 Mech Teck 7. 622 62 65 7 8/14 2 4 (177) 6 5/1 591 4 (177) 6 5/1 591 4 (177) 6 5/1 5/1 **23/H \$/**4/92 6/22/9 1/0/92 @ 5/66 92 8/3/2 (A) ~ <u>IJUR 6/22/92</u> - Stilling Equipment/Component Identifiers Serial Number Room # Gloveboz Manufacturer Measurement List Parameter Name Model Number Symbol LANL Property Humber Itom Description Other Equipment/Components To Be Identified? IN/A Form No. WELF OF Revision No. 3 Data December 12, 1990 Page No. 1

ć.,

24

1

e,

WIDR (continued) Commente widr # CONDITIONS Stack failed No voltage of the steek is failed because of low flow then لاسمانه سما o measure the EM but set a flag industing - 2 Actions the low flow condition MEASTER RM x × if the stack is failed because of led instriment "4ail" . X power years ** t** don't measure the EM and store the failed × condition in the measurement value. Action Taken added the ability to set a low flow flag in the stack report file when The stark ATT is second during low flow conditions. added the ability to add an "f" to Res values measured during low flow conditions when printing the steek reports. Action Completed (TSM) **Completion Date** Form No. WETF 01 Revielon No. 3 Dete December 12, 1990 Page 2

	Test No. 448				
	Paget of				
as specified w a) low flow co b) no instrume	hen: ondition if gow		courts fanation	Preliminary	By: <u>O. Hart</u> Submitter Reviewed: <u>Al Mayula</u>
c) normal oper	، مرسط ص	ord trans	· · ·	Responds to WIDR (s)	Approved:
		TEST PREREQU	JISITES		
Test S/W Installed	Ver.	Configuration Used for Test	Verified By	Facility Configuration	Verified By
ETC.	2.27	Normal		No Alterations Required	
		• • • • •	 Test Conductor	Alterations Required. Ref.	Test Witness
	I	TEST RESULT	S	Reviewed by	and Date
Test Results are a Retest required.*	cceptable Referenc	e Test Log for:		D. Ha. +- Test Conductor	<u>Helling R. 7</u> Test Witness
		POST TEST AC	TIONS		
Op. S/W installed	Ver.	Other Actions	Verilled By and Date	Facility Configuration	Verified By and Date
yes	2.27	u/A	O. Hand- Test Conductor	Restored to Pretest Configuration Image: Configuration As Noted below Image: Configuration	154 Test Witness
a/3/92	To-			J NIA	1

n.

4

5

긝

. .

1		Page 2-01 2-
	While Fuc/arthe truning, bring up the BTE Z.Z. Note time at which cade erres up.	Conce the
	Les vie eperater per a minimum of two minutes the in two with the in the contract of the flaw. Set stack flow alarm wing 20031 (EFH-Exitist, 2143). This will generate a family demande a family demanded a family	te e laur 41
_	ararmy as well as the TH STREE FHAL THAN THE WE ARE APARTE for a minimum of a minimum of the more share at which the low films was initiated.	
ю Д	Effect strack plan alorm using 20031. Researt strack monther fail using 21100. Note times 2t which performed, hest orce operates for a minimum of 2 minutes.	Note times
-√1 > -√	Set had voltage alorm using 20031 (EVA-EXHSK) DI A). This will generate a "TH- STK Volt Status" alore se und as the "TH street FRI: FLRG". Ind one should be a	: [*] тМ-`5тК \ a
	Minimum of 2 minuter, noting the mine of which the book was initiated.	sted.
4. 87.5	Reset voltage bad status aloren wing zeosli Res markeur wonder wing zilde. Note times at which performed. Let the operate for a minimum of 2 minuter.	.106.
<u>ي</u> ٦ د	Using 21560, 34mg the BTC contro. Note thing . 24. Which code shapped.	.
	M	" uthere
÷ ¥ Ř	filename markely the stack date filename, e.g., struck_120131. DHT. The proof report is printed on the printer. Verify that stack report reflects actrons at time taken.	lects .
		·
	•	

÷,

 7/31
 10:48:34
 HVM-H312A
 RACK- H3 +12V
 PWR, ANALOB
 9.42 L0

 7/31
 10:50:14
 EZS-ETCRUN
 ETC SMODE RUN
 1
 0K

 7/31
 10:50:24
 GGH-LI
 LOAD IN- G.B. 02
 15.68
 11

 7/31
 10:50:24
 GGM-LI
 LOAD IN- G.B. 02
 15.68
 11

 7/31
 10:50:28
 GGM-LI
 LOAD IN- G.B. 02
 15.88
 11

 7/31
 10:50:28
 GOM-LI
 LOAD IN- G.B. 02
 .16
 0K

 7/31
 10:50:28
 GOM-LI
 LOAD IN- G.B. 02
 .16
 0K

 7/31
 10:53:23
 GOM-LI
 LOAD IN- G.B. 02
 .16
 0K

 7/31
 10:53:23
 FLM-WH20TK
 WASTE WATER TANK LEVEL
 1148.23
 0K

 7/31
 10:53:23
 HVM-H312A
 RACK- H3 +12V
 PWR, ANALOG
 12.19
 0K

 7/31
 10:53:23
 HVM-H312A
 RACK- H3 +12V
 PWR, ANALOG
 12.19
 0K

 10:55:32
 10:55:32
 10:55:32
 10:55:42
 12.19
 0K
 12.19
 0K

</tabula U C A C 4 C A A U 10:55:32 OTID 1125 - DIIDS - CHANGED TO: SET 826.35 HI 7/31 10:55:43 WPM-TK5L1 SLP1- PR. A 7/31 10:56: 6 EZS-ETCRUN ETC SMODE RUN 0 LO 10:58:27 DIID 1104 - DIIDS - CHANGED TO: SET 10:58:39 43 - DIIDS - CHANGED TO: SET DIID 1 HI ' 10:57:35 DIID 43 - DIIDS ~ CHANGED TO: RESET 7/31 10:59:36 EFA-EXHSK _____TM- STK FLOW ALARM <u>o ak</u> 1 ŰK 7/31 11: 0:10 EZG-ETCRUN ETC SMODE RUN 676 Cola Storting 11: 0:17 DIID 1104 - DIIDS - CHANGED TO: SET 7/31 31: 0:20 EZS-ETCNULL ETC CHODE NULL ETC IN FUTO ***** THE TIME IS 11:00 THE DATE IS 07/31/92 ***** 0 040 11: 1:30 'IID 43 - DIIDS - CHANGED TO: SET 1/31 11: 1:31 EFA-EXHSK TH- STK FLOW ALARM initiate low plaw 1 HI Δ 7/31 11: 1:33 EZS-EXHSKFAI TH STACK FAIL FLAG 1 HI 61 11: 5:30 43 - DIIDS - CHANGED TO: RESET DIID 7/31 11: 5:31 EFA-EXHSK TH- STK FLOW ALARM reset low flow 0 OK A 11: 5:40

 7/31
 11: 5:40
 EZS-EXHSKFAI
 TM STACK FAIL FLAG
 #547
 \$129
 0 0K

 7/31
 11: 5:40
 EZS-EXHSKFAI
 TM STACK FAIL FLAG
 #547
 \$29
 \$129
 0 0K

 7/31
 11: 5:13
 HVM-H312A
 RACK- H3 +12V PWR, ANALOG
 9.52 L0

 7/31
 11: 6:13
 HVM-H312A
 RACK- H3 +12V PWR, ANALOG
 9.52 L0

 7/31
 11: 6:42
 HVM-H312A
 RACK- H3 +12V PWR, ANALOG
 9.52 L0

 7/31
 11: 6:42
 HVM-H312A
 RACK- H3 +12V PWR, ANALOG
 12.19 0K

 7/31
 11: 6:42
 HVM-H312A
 RACK- H3 +12V PWR, ANALOG
 12.19 0K

 11:
 8:36
 8:36
 11: 8:36
 12.19 0K

 DIID 1144 - DIIDS - CHANGED TO: RESET U A. 11 A 15 11: 8:36

 4 - 01103
 CHANGED TO: RESET

 7/31
 11: 8:36 EVA-EXHSK
 TH- STK VOLT STATUS () had bed forer supply

 7/31
 11: 8:39 EZS-EXHSKFAI
 TH STACK FAIL FLAG
 1 HI

 7/31
 11: 9: 2 FLM-WH2DTK
 WASTE WATER TANK LEVEL
 1201.58 HI

 7/31
 11:10:41 WPM-TK5L1
 SLP1- PR.
 950.06 HI

 7/31
 11:13:21 FLM-WH2OTK
 WASTE WATER TANK LEVEL
 1146.13 OK

 DIID A 11 A A 11:16:31 4 - DIIDS ~ CHANGED TO: SET DIID 7/31 11:16:32 EVA-EXHSK TH- STK VOLT STATUS TASET bad for supply ok 11:16:40 DIID 1144 - DIIDS - CHANGED TO: RESET 100 + 101 + 109 0 OK It 7/31 11:16:41 EZS-EXHSKFAI TM_STACK FAIL FLAG .1:17:31 JID 1125 - DIIDS - CHANGED TO: SET

 7/31
 11:17:32
 EZS-ETCNULL
 ETC CHODE NULL
 ETC Code STOP
 1 HI

 7/31
 11:18:4
 EZS-ETCNULL
 ETC SMODE RUN
 ETC Code down
 0 L0

 7/31
 11:20:53
 WPM-TKSU2
 LPR PRESSURE (U2)
 351./1 HI

 7/31
 11:21:1
 WPM-MB4
 DRY22- TANK PRESSURE
 529.91 L0

 7/31
 11:21:3
 WOM-HX60T
 MO-2 02
 MONITOR
 .27 L0

 A C E A A

STACK INTEGRATED T2 REPORT FOR 07/31/92

hoursesses

increase in stack integrated 12 during the minute. In mCi

	T 0814	
24 Wr	accus	oiss
Long-1	term acc	

anddyy hhamas oCi last reset at: miss miss pCi last reset at: miss miss

· ·

06xx 07xx ain. 40 00zz 01xx 02xx 03xx 04xx 05xx 05xx 09xx 10xx 11xx 12xx 13xx 14xx 15xx 16xx 17xx 10xx 19xx 20xx 21xx 22xx 2.74 miss als: niss. miss **a**|55 aiss Miss mtes mise mise miss miss miss miss nisa =ias **mi**28 mise mi sa a i e e **m**[88 nias o jes miss miss 01 02 M155 míst miss តារីនដ miss RÌVS niss **a**l 98 olsa miss mias miss 'mlas ตาธร niss niss miss aise **n** 99 aise nšes lof miss mi ta miaa **mi** as mise m aa miss aiss miss miss miss 🍳 miss miss miss aíss **B**¹53 **B**İ 55 ดาระ aiss niss mist piss. **miss** mios miss miss miss 3 mias miss 9 03 04 05 06 mise **8** ni es 10f miss Bitt. mtan. **ai** 35 mi sa ntes ntes mina mina misa mias. mi as สารจ aiss mias. ងរ៍នង **mi**88 mias ៣វិទន mins miss piss ntsa of miss miss miss m 58 0195 0151 mi 68. ៣ខែន mi 55 mise mi se aiss. aíse **Q**[31 miss atian. **#1**68 ແກ່ອອ miss miss aise Bİ 55 misa Of miss ains ales aiss BÌ SE alse MISE M 96 miss. nisı niss aiss alss al ce Dies. **B**| 85 mias miss mlaa miss miss miss miss miss' miss miss miss MİSE. wist. mise miss. atin. miss mige. niss _____ 07 alss. mi 55 mi ss MISE ai sa miss. ពនៃន mêaa miss miss mise 0 alus miss fieil miss miss miss miss miss inist. miss **1168** miae mise miae miae nies miee 08 09 aise miaa miaa miss mias **m** 1 8 8 mi as misa niss. mi sa ារំទន **B**125 miss **nise** mias 0185 ats ats miss misa misa fail misa misa fail misa misa fail misa misa fail misa misa fail **ai 63** bi 55 ni ss **ai 3**8 mi 25 ntee mtee at as **mi 98** mi 65 m198 miss mlea miss. **#1**89 al 69 mi sa miss **D** 88 **mi 3**6 10 miss miss al ss R 35 miss mias **mi 66** mise mise miss **B** 89 0 88 el sa ntss **B**|55 nisı. mi sz. al ss **#1**98 mise ates ai ss mi 13 mi ss ៣ខែទ miss **B1**35 mi as **el 9**8 **e**| 38 mise mias mise m| 85 MISE miss ៣រំទន uise. mi ss mise. **#**388 misa 11213141516 miss miss piss m| 65 alss aiss alss mi se. 88 **m**| 69 **1**88 **mi 6**8 ni sa mi es **a** | 64 **69 10** nisa ni se mitt mi ss alsa mise mise alas mise mise alas mi se mi ss miss aise. miaa miaa mine **bise** mi sa aisə **mi 96** mies mise aise **B**| \$3 ni se ni sa miss miss uige miss miss uige miss miss uige DISS AFSS AISS **a**| 68 pel se **ni 5**5 **mi 65 B**| 55 m188 **m**| 96 missőifeil tal 88 Mİ 69 **n**ias **B**| 66 Mİ 96 m658 MISS IFAIL mias miss miss miss miss miss miss **23 1 m**i 65 **ni** 88 miss M 88 **B m\$86** alss. niss aiss aiss e î sa misa mtea miaa 165 n en **ni 3**8 misz misz int se 83 10 mi 85 miss. ni se mite M 85 **m[ss #**\$98 mina 17 mine mins mies Ri 45 m! \$5 mise el es misa mias mise miss miss miss miss mi su miss mias miss m188 mí as miaa efae miss â aiss aiss 18 atse. **m**| 93 mtes mise mina misa misa **ni 5**8 **B**| \$8 miss mtss. niss. mi en miss ii i sa misa mfss. #1 65 MISS MISS MISS MISS MISS MISS **B| 55** mtso mias miss 19 20 21 លខែទ ពារំទន ais: **ní**sa **ां** हत्र aist mlaa miss. **mi** 88 mias mise. mles. mtas **A** | 88 miss **M1**58 **B**| 88 mi 68 miss. miss mias. mise . aiss ៣វែង Mi 81 **A**| 15 m155 BÍ \$3 miss miss miss nies mi 10 miss mise ∎tsa mi sa misa miae miee afse mi sa aica aise aist ∎i 98. រារ៍ខង miae. ៣ខែង miss alss miss misa miss. ari sa piss mi sa mi sa mies mias misp. **R**Í 55 _____s ∎i es 2242222829555555555555 mi sa mi 15 5 niss mi as mias mi 68 mi 55 mian mlsa miss miss 0155 0155 0155 0165 A | 13 A | 18 **miss** miss miss miss miss miss miss miss alss CO 55 miss nias. misa mise **155 135** aite aist. mise #\$56 ai es misa misa misa Nisa misa misa Nisa misa misa Nisa misa misa Nisa misa misa aise niss. miss atsa misa niss mias miss miss miss **mis**# **#16**9 ៣វីមន **M**]88 **aits** aíst miss **al 8**5 mi es mias A153 alse. สายม @158 m199 miss mEse miss Miss mias. mise. **#1**88 aeim mi 88 aeim mf88 ∎f s\$ miss #155 .#164 **1**169 wias ats: misc miss. mi sa **m**[98 aiss mias. eí sa aiss. miss ariss niss **mi 95** ៣ខែង ៣ខ៍ខន alse miee **#**[28 miea misa asim. mi es oni ee 01 6 S ntas atas គារិតន miss mi 88 miss mi au mi sa ណ៍ទទ aisı #18# mi ss BIAS BIES mito mice mise **#[68** mias mias n s . ales mias miss nî sa mî sa ales ៣ថៃទ mine miss pise. **N**| 35 miss mise **185** mias miss miss miss #1 5 S #1 3 S miss miaa miss miss at az mi 35 sa is miss mias mias miss mise miss. លាំនង **m13**5 miss atse. **m**i 80 **=|** \$\$ mine mise mise mise mise mise mise mise mise **189 m** | 63 miss **1 85** ari se miss miss at as nios nias miae etsa al sa **m**| 98 mles #196. mi 65 ណរែច mias **B**| 55 m| 56 ៣ខែទ mise mise miss **m**| 65 nies miss miss 11 22 aiss. aiss. mi es miss. miss mise mi ee **m(**88 el se m688 miss MISS miae **m**| 88 **B**165 al si **1**89 míse mi 99 miss niae mise atas - ed 68 miss eiss. miss niss afs: ni sa mi sa pies miss miss ml 88 miaa misa mine at as mise mise mise mise miae aeta mias nisa **m**| 55 RÍ 55 m168. n|5\$ mius p| \$5 ៣វិតម **=**158 mise miss nies **B| 15 a**| 35 **m**] 24 mias mias atsa · ni ss miss mi B-S ni sa **al 68** ai 35 លា៍ ភន mi 65 mise m(\$8 miss eiss aiss **a**i 33 miss miss miss misa mi 58 miae mtee e las **ni** 69 miss miss mi 65 mt 69 atss este m‡58 **61 15** mi sa mist **m**l 68 , mtaa **ni** 76 ៣វិនគ m185 miss miss miss miss miss miss plas miss miss miss plas plas oiss siss miss miss **1145** otes m185 **m**| 96 mi ea mi 89 BÍ 55 **B**Î **S**S miss miss miss. **e**188 miss miss miss miss miss miss **B| 44** aics m[33 miss . ៣វ៩៩ **M**| 66 miss aiss 2111 mi 29 M158 mi es misa កាត់ទទ **Gise m588** mi as mtaa m108 m128 mise mise mise mise mi sa mite mite pice also mise also pice also miss miss miss **ni** 69 miss miss miaa mi se m\$98 miss. mias miss al se alss MISE. កាំងង miss ារ់ទដ mise m| 88 **m** 38 mian mise RÍ 55 mi as miss miss M 95 ធាំទទ ml 69 mtee miss miss miss miss niss ni as mins **mi 23** alsa mine mies mies mies ntes ៣នៃទ ates mt s-s 10 65 miss **B**| 55 stss. af Ba miss mi 22 mi 55 mias mbaa #168 ulse mise mias mf## ani ne niog miss miss piss miss miss miae **8**8 mise aise. ni sa mi 65 m1 88 miss mlas mi ee **D**ise mi 88 miss ae im **28 10** mius ni sa ៣វៃទាន mi -655 atas mias mias mias mi 96 mi ca miss niss miss Miss miss miss mics miss **=1**66 mise mise miss mise mies aiss ∎Ì₿₿ misa miss mies **- 86** miss mise. **B**| 15 **N**| 55 mtss mtss. **88 I** ai se mica លវត្ត miss miee mise **mi** 68 mtes mt sa pies. mise mise =Est mine nine nies nies mias neim salm R| 59 miss miss mias nies alss aiss mias miss miss miss aiss miss miss mies mies miss miss ៣នៃទ mias niss mi ss aiss miss miss mlas m(69) m i é é niss miss miss miss alle alles alles after ത്ടം miss miss miss atas atas. mi sa នានៃន mi 53. **m**i 98 **M** 88 #8in miss. **3** 89 misa ni es mise miss miss miss mise miss miss miss **28 1 mi 6**8 mies m188 mi 68 mias **e**[55 miss as im anias ពានៃទ aiss miss miss miss ពារំនង miss. niss **ni 5**8 2310 miss. **mi** \$8 @156 miss mise **m\$80** mise miss miss m\$88 **B**158 B 85 mi#3 aiss mí se mias. miss miss asta miss miss mi 15 ៣ខែង mi 55 m155 ការឧទ mise mise miss **M**135 miss mise wiss miss miss mise miss. **M**¹63 #\$\$\$ mi 69 mies mi 63. mies mies mics miss miss miss mice miss miss miss **m**] 86 ៣ខែទ miss mise ារទេ mise wies mice miss miss. miss. m188. **...**198 **B** | **B B** mi se mi 69 mi 88 **B**| \$8 mise ae in mi #8 **#166** mine **#158** mias aiss **... mi 65** mias atas mi es misa tain t ៣ខែង **aise** ei sa mise. mise 81 18 Õ mi 15 mi 19 MISS MISS mias **169 0195** miss m198 mi 55 **F**| 96 MÎ 65 ៣ខែទ mias mise ៣ខែន ៣វែនផ mias miss miss mias. **mi**## al se mt sa misp miss mise nise aias mi es Ó mise mias miss mise mias miss mise alse aise al co niss. mine 21 I A 22 I M m| 63 Õ M| 85 A| 55 miss miss mias miss m199 alss tal 88. **1** 2 2 mfaa miaa mi sa **a**|35 es fa ntes miss miss mise miss al se mice **11** 2 2 mias. m (sa niss miss mlsa លវៃទទ mina mi 85. niss Diss mi \$\$ miss mies mies ៣ខែច **ni** 52 mişs mias mi an m1 89 miss aiss **aise** miss B| 35 **mi 55** mi 68 ei sa **mi** 38 eise mise @i68 mi 29 miss miss mi 55 mi se miss **B** | 65 mi ss mias สารธ atin zeim mi 66 mi 56 **m**|88 pal es aiss M| 86 M| 55 **#1**99 #[96 miss mi 68 mi 99 miss stes mise aise Mi 88 M 55 niss Alss R 81 mise nica mi os miss **B**| \$\$ M 89 8195 miss **al** 86 ៣វិទន miss mi ss **#155 m286** niss aa la astm asim astm mi sa miss mi 35 p186 mise mise **m** 28 miae miss aiss M 68 **mi s**s M 55 mins aiss misy misy alsa miza siza miza misa mico miaa nies mias **B**|55 miss m183

A A ALL AND AND A REAL AND A

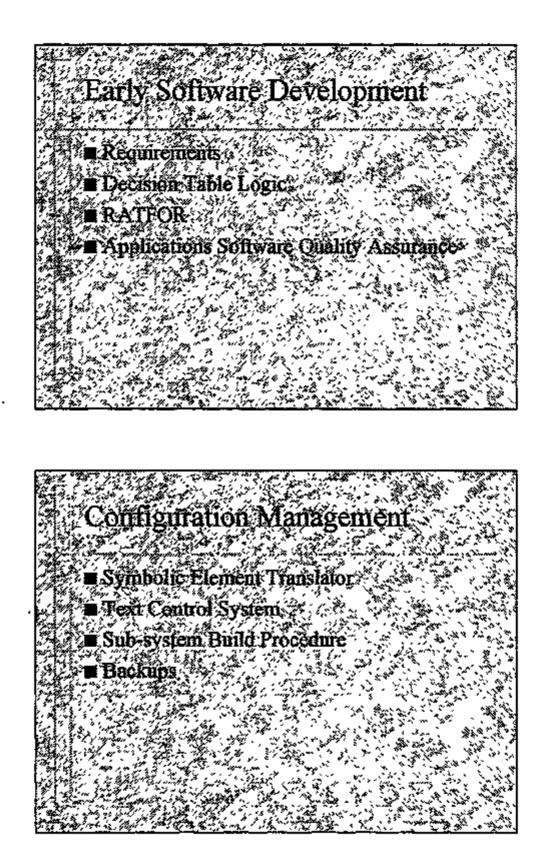
.

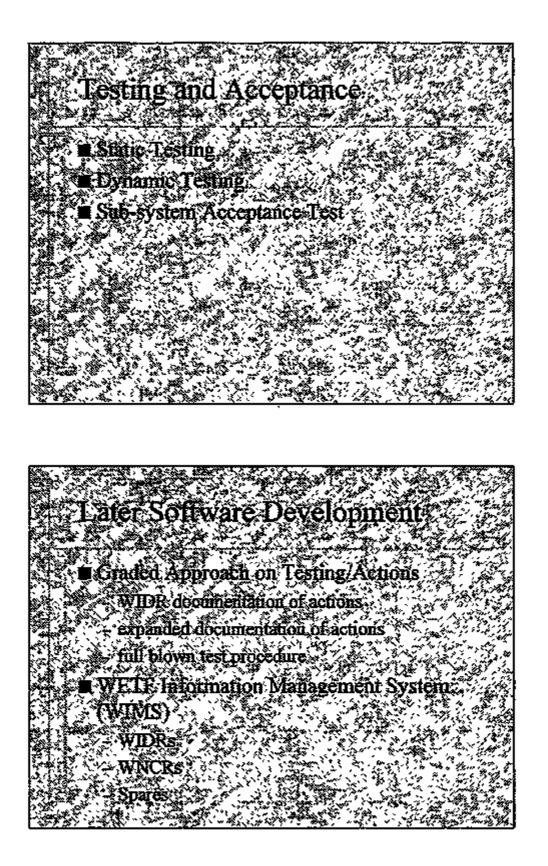
. . .

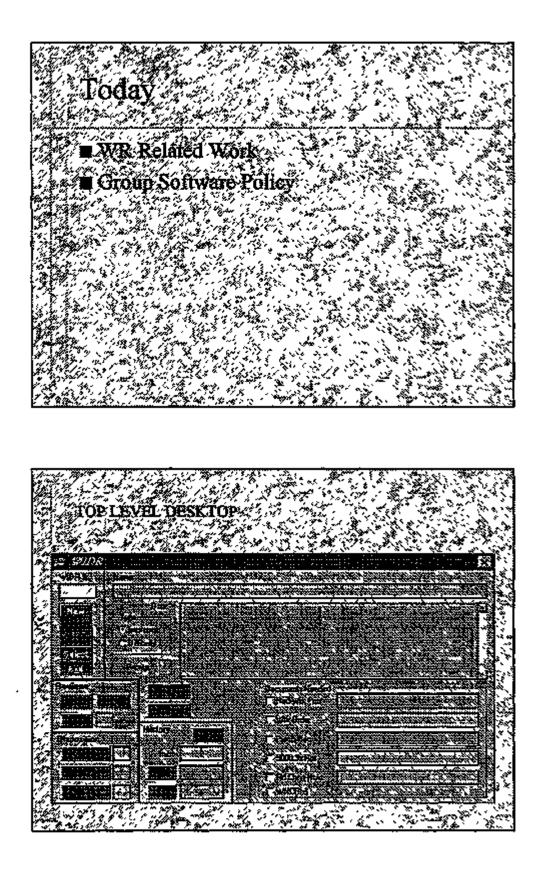
. -س

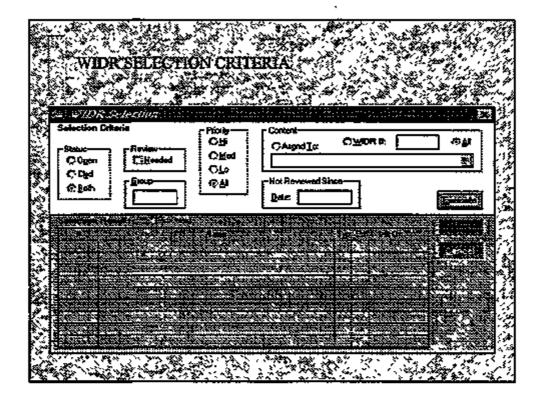
•

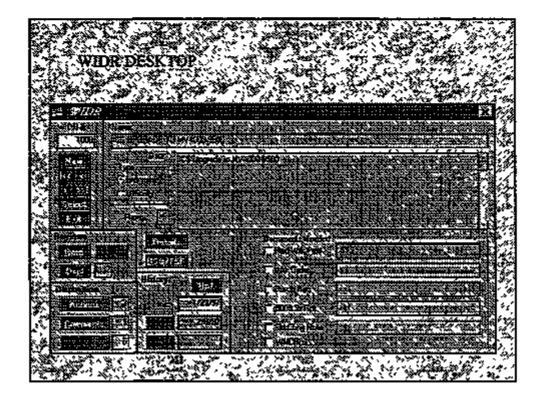
I











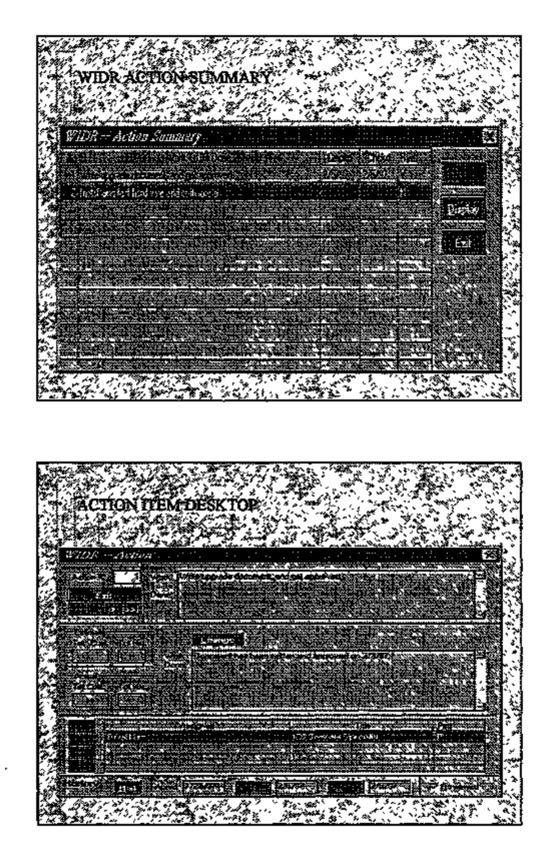
I

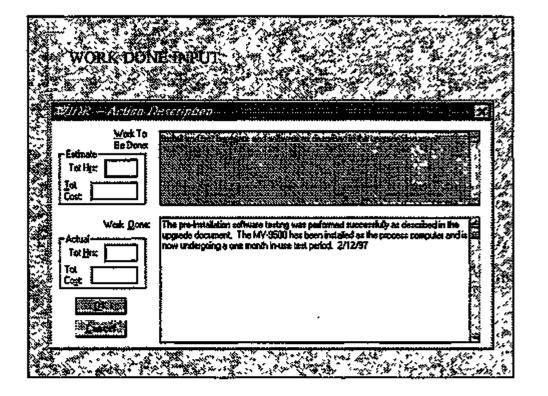
| | |

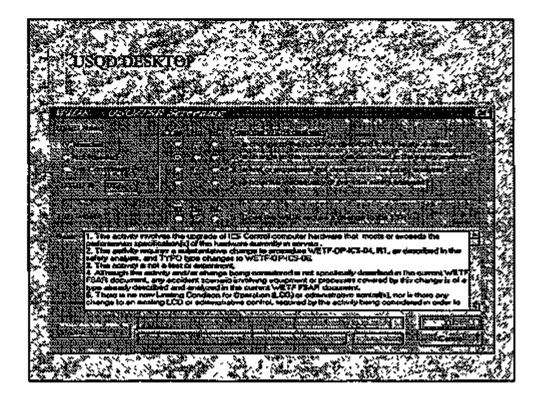
ł

1

ł







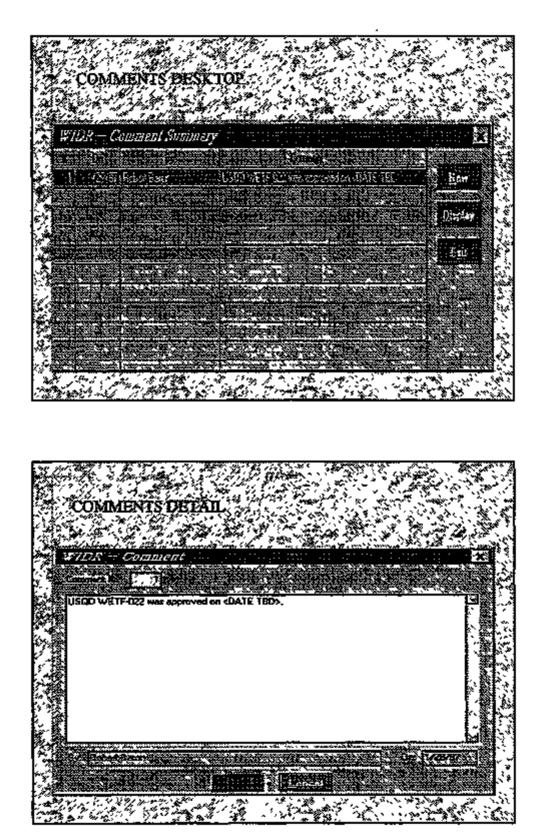
1

L

i

)

Т



•

•

{

C;Puch Heg	Coourseels Needed C, Rud-Witj Pint	Ling the second	CISAR @Incoverent Ciplinus	стор — <u>ж</u> ини. IC3 UPERADE T0: MV4000.9000	WDR DETAILS
	j Soon Jr Ber	131 näym Cankeninskinn 152 Cogysienser 131 näym Aljanna 131 näym Aljanna 131 näym Aljanna	Area		

•

()

.

							26° 05 200	2
			2/2/14-4	70 1 4000	324	老い.		έ.,
	1 - S			Sof kizi -		20	**** ×	١÷
擾	n se			21840			25% H	T)
						> 2	<u>کر ک</u>	í.
6353			822.112		33122		_*″ ⊴ <u>∏</u>	Ľ
78.22	200000000 2000						79. S.	X
1712	REA				1	8	and the second	R
14 18				4148	1853		ଞ୍ଚାନ	Ś
1128		御に 森	惊胁				2	í)
叫儀				o HB	2012		9. sof	2
関係							27 J	ľ
H					636		í . 'S s	L.
		t: 1 🖉 🔊	14.921		89392		N 8. s	4
粗膠				809.4	19957			٤,
毗露							\$ V. S	
88	202		1362		192320		1 A S	ž
目機					12686		ಗಳು 🍼	٠Ļ
月識		KI MA					1994 / S	đ
11 22		IMZR				££	E G Ø	3
目缀	劉麗伊						3 . M	14
11 🏼					10/10	5 (C)	Se fi i i i i	2
li XX			译试出			- S	Set .	1
11 122	劉朝朝		84 2 [NA 3	٩â
1422								2
						1. S	$\Sigma \Sigma \delta$	*
						an i∌	1.1	1
11.	22 M				123		<u> </u>	3
162	628			8 <i>7</i> 872	312			¢,
1180							¥ X	1
							143	22
	8121 3					- S.	1. S. C.	.,9
1.055		123183633		31531231 2433			8 yr - 1	κş
100000	6.777.7 <i>8</i> 6			SS 332773	365X 138		S. 31	

Orval Hart

۱. ا

1

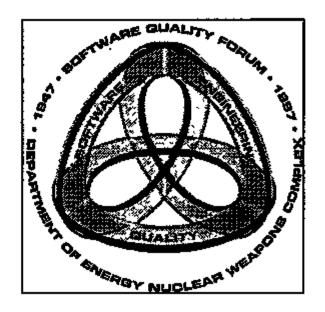
1

٩

.

10

 $\overline{\bigcirc}$



Session A2: Software Engineering Processes

Chair Kathleen Canal

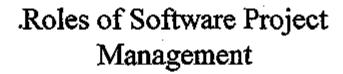
Headquarters Department of Energy

Session : Paper #	Author(s)	Title
A2:1	Michael Bell Lockheed Martin Energy Systems	Function Point Count Adjustment by Means of Scaling Touched Function Points
A2:2	Stewart Meyer Westinghouse Savannah River Co.	Using An Automated Code Management System To Improve Configuration Control Practices
A2:3	Karen Jefferson, Terry Porter & Todd West, Sandia National Laboratories	Software Engineering and Graphical Programming Languages

Function Point Count Adjustment by Means of Scaling Touched Function Points

Michael A. Bell

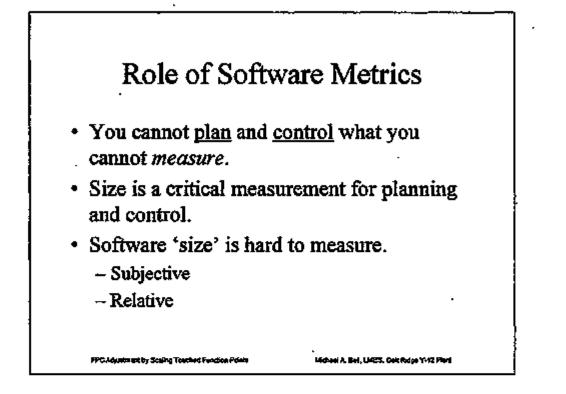
Lockheed Martin Energy Systems, Inc. Data Systems Research and Development Software Engineering Oak Ridge Y-12 Plant

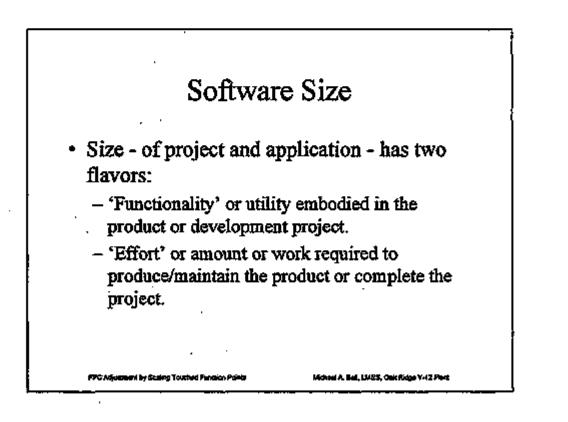


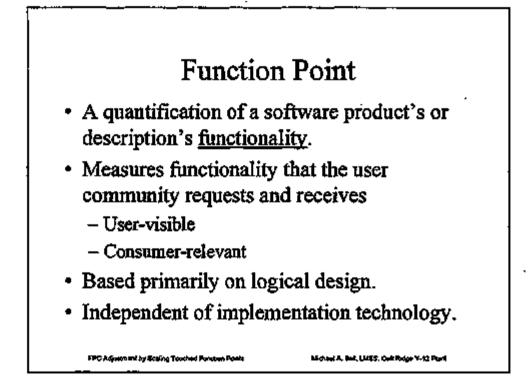
- Supports software development projects and application support.
- Plan
- Control

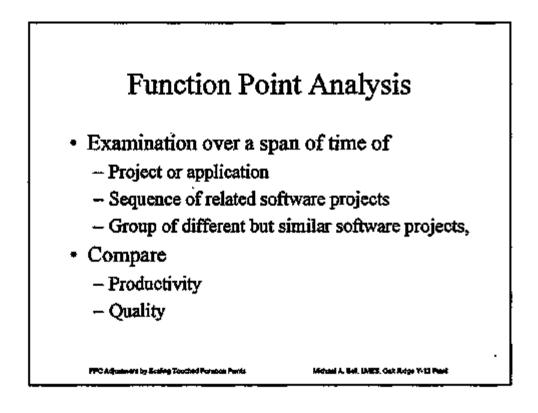
teant by Sailing Teached Paralies Points

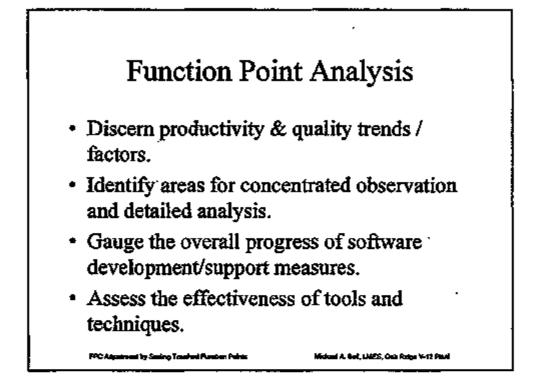
Michael A. Bol, Links, Oak Ridge T-12 Part

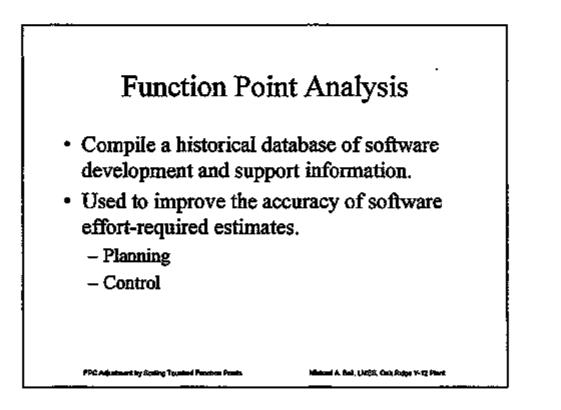










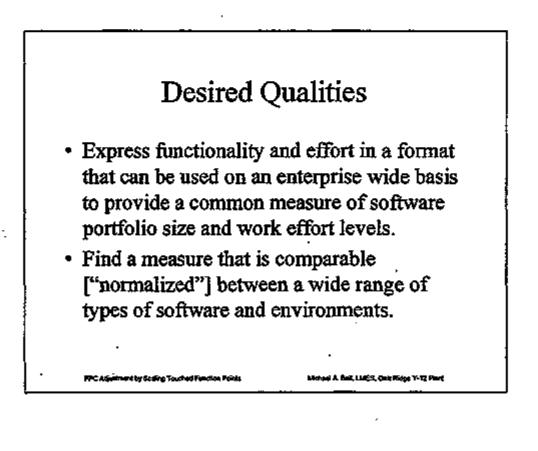


Two Uses for Function Points

- Express the amount of <u>functionality</u> delivered or supported by a given effort, independent of the technology and implementation details.
- Estimate or express the amount of <u>effort</u> required by a given effort, *taking into* account the technology and implementation details.

el A. Bell, UNES, Ock River Y-12 Prov

n eni by Scalaro Touched Function Po

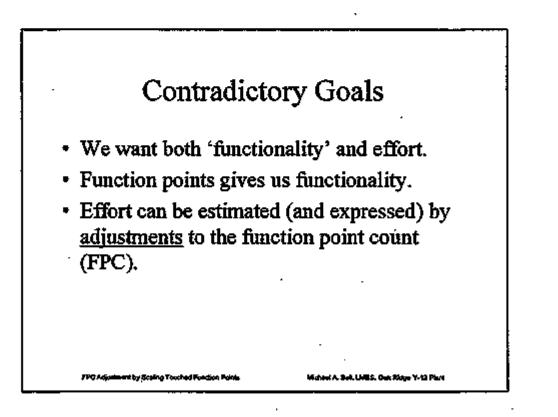


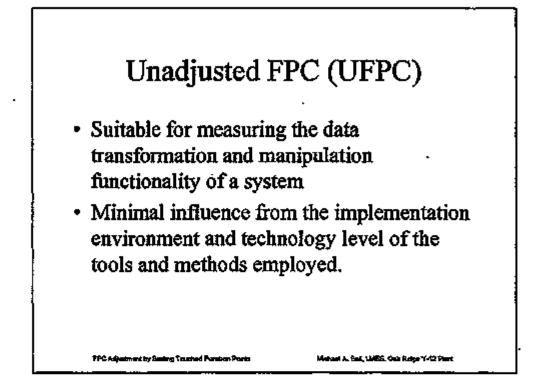
\$

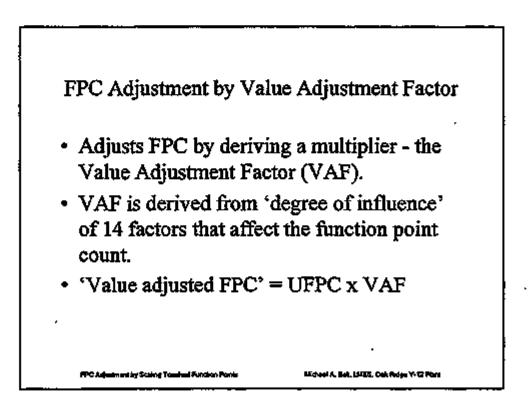
Contradictory Goals

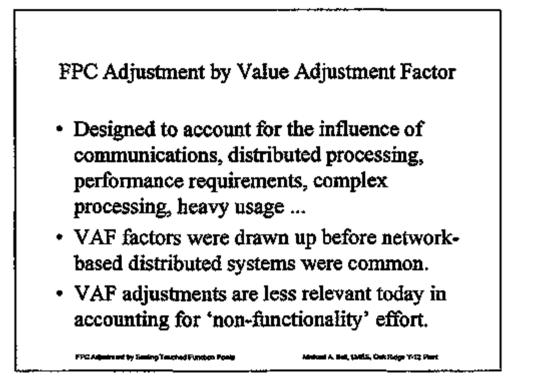
- Factoring in technology differences and other factors reduces the degree a function point count measures "pure" functionality.
- Leaving out technology differences and other factors reduces the degree a function point count measures actual work effort required for a given project.

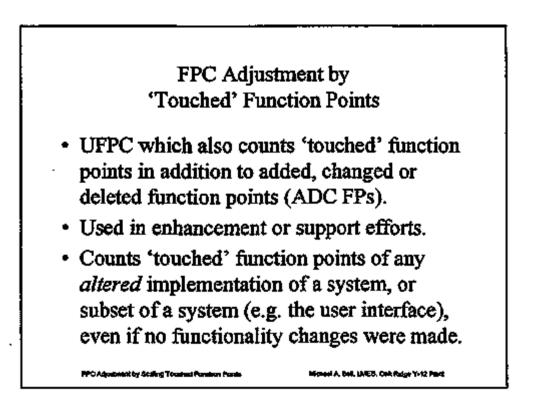
A. Dell LARS, Oak Ridge Y-12 Ph

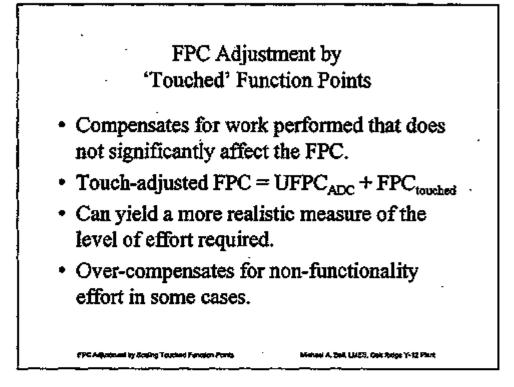


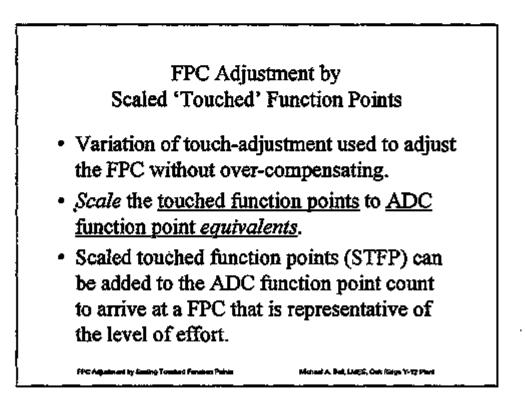


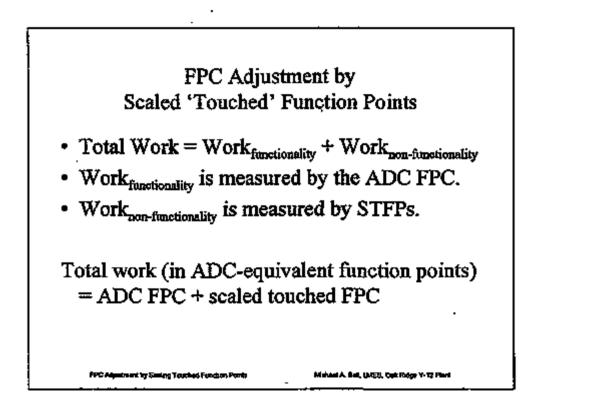


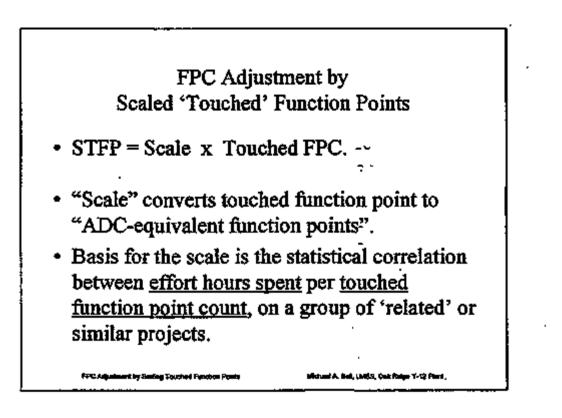


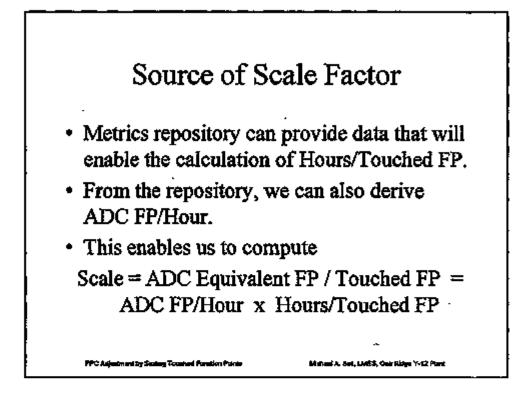


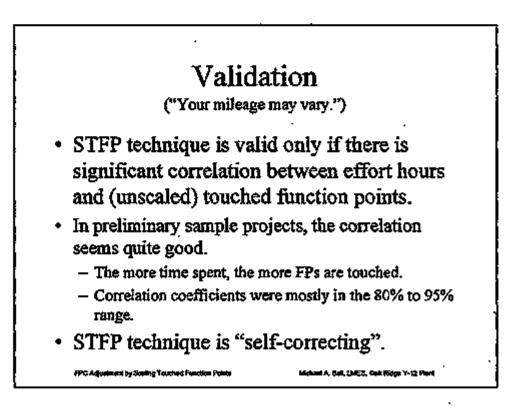


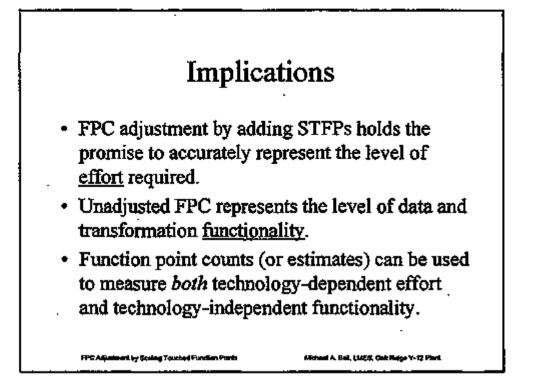


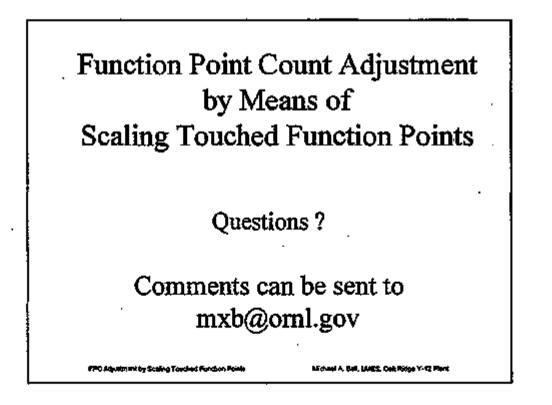








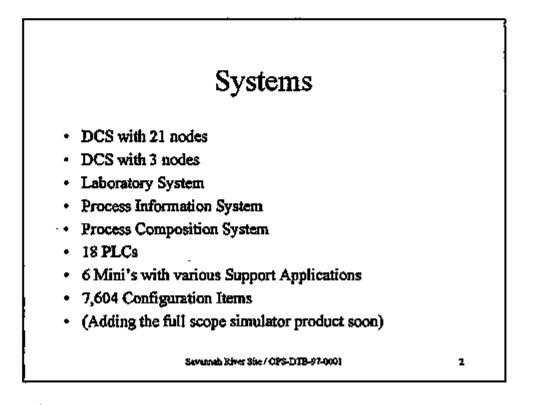




USING AN AUTOMATED CODE MANAGEMENT SYSTEM TO IMPROVE CONFIGURATION CONTROL PRACTICES

Presenter: Stewart Meyer

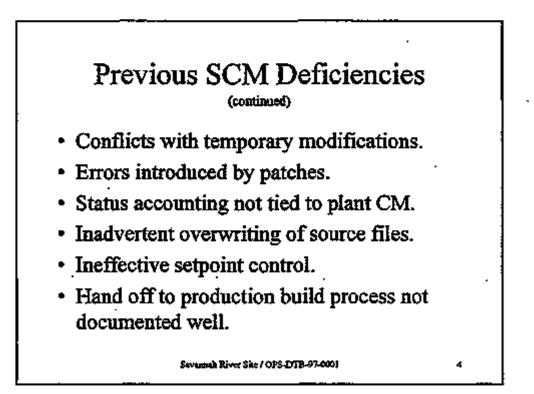
Sevenah Biver Site / OPS-DTB-97-0001

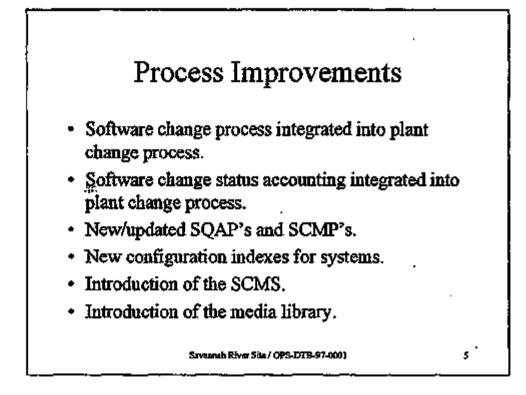


Previous CM Deficiencies

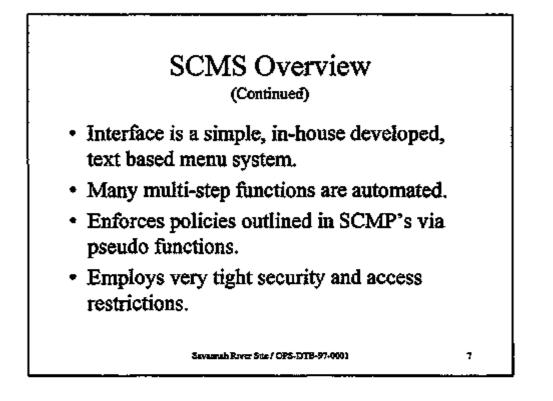
- Software documentation not integrated into plant CM process.
- Used a directory hierarchy for development vs. baseline.
- Change sets were entirely in paper.
- · System backup the only protection.
- No audit trail on modules.

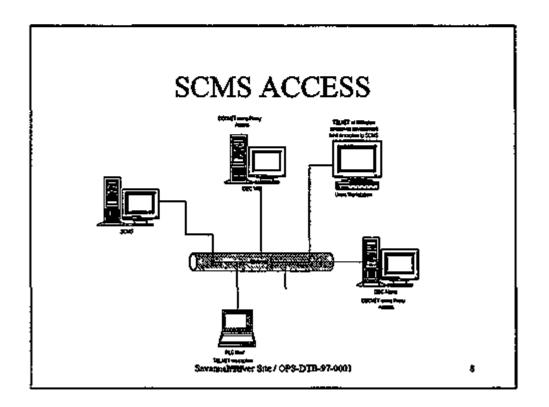
Savanah River Site / OPS-DTB-97-0001





<section-header><section-header><text><list-item><list-item><list-item><list-item><list-item><table-row><table-row></table-row></table-row></table-row></table-row></table-row></table-row></table-row></table-row></table-row></table-row></table-row></table-row><list-item><list-item><list-item><list-item><list-item><list-item><list-item><list-item><list-item><list-item>



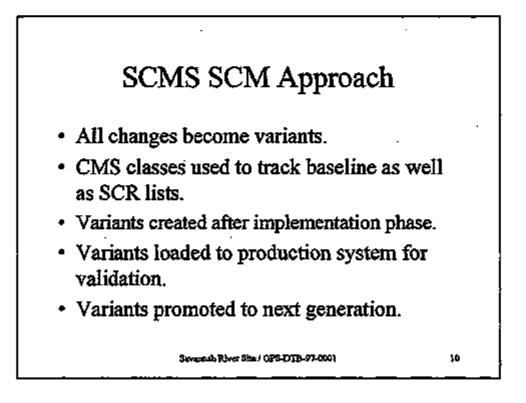


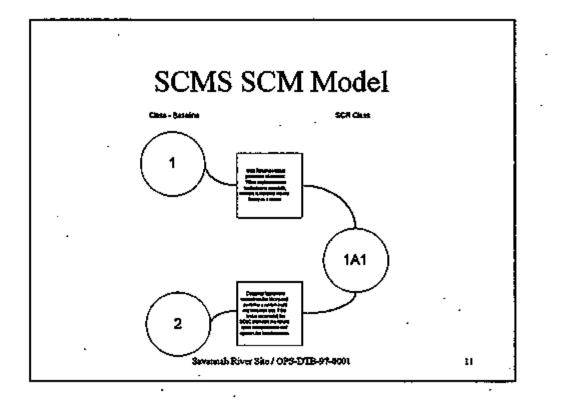
SCMS Security

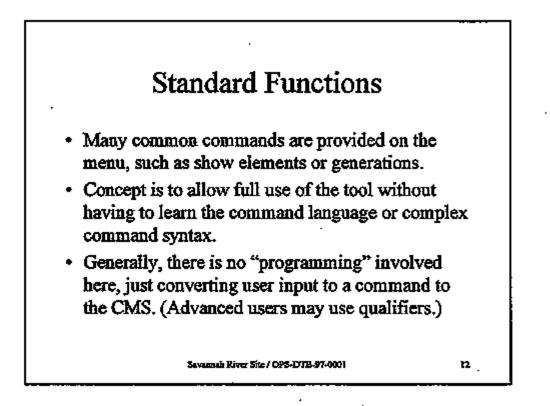
- Uses either Proxy or Captive accounts.
- General users cannot perform tasks at command prompt level.
- Access control at the OpenVMS level supported by additional ACLs at the CMS level.
- Several levels of access enforced.

Savamah River Siz / OPS-DTB-97-000)

9







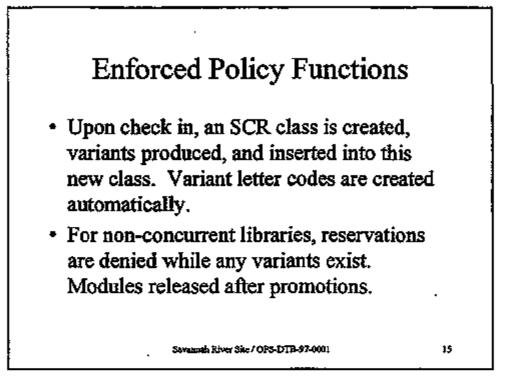
Enhanced Functions

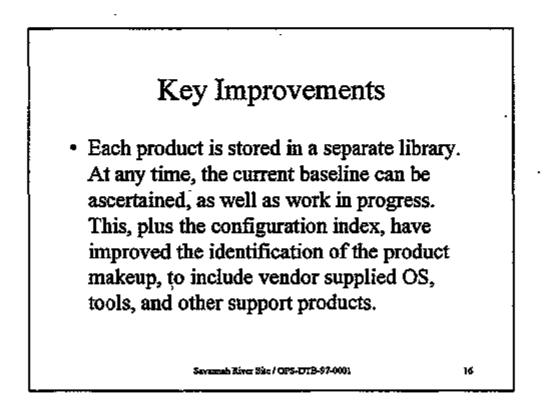
- Check out by element, group, or class.
- Check in by SCR number.
- Different library history types.
- Class merges.
- Automatic merge class creation.
- · Management reports.
- Transaction comments generated.

Savagnah River Site / OPS-DTB-97-0001

13

Enhanced Functions (continued)
FTP file transfers to workstations.
Promotions by class
User log file.
Empty and delete a class
Saved user configuration.

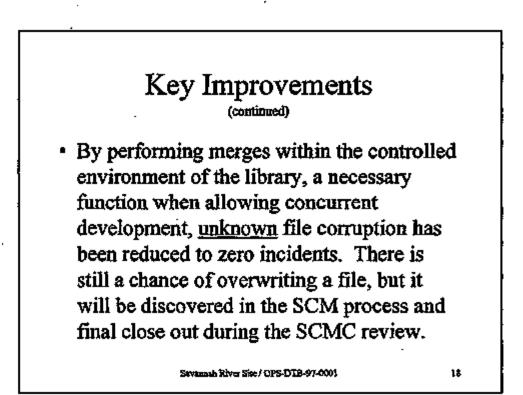




Key Improvements

 By using the SCMS functions provided, patches are now stored and verified.
 Reports from the developer produced after verification are checked against the same report run on the production system after the patch is installed. This provides instant feedback on possible errors introduced due to typo's or incorrect field modifications.

Savannah River Site / OPS-DTB-97-0001



Key Improvements

 In using the SCMS we can now perform periodic verifications on controlled systems with confidence. Executables as well as source may be subject to this control and review. The elements in the library are compared to the equivalent on the production system.

Savannah River Site / OPS-DTB-97-0001

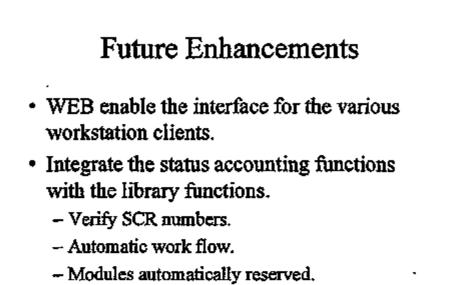


• Configuration audits are now much easier. Using the group or class contents we can produce reports on the current status of any library. There is also a separate status accounting database application that, when used along with library reports, provides a clear picture of product status, schedule implications, and resource assignments.

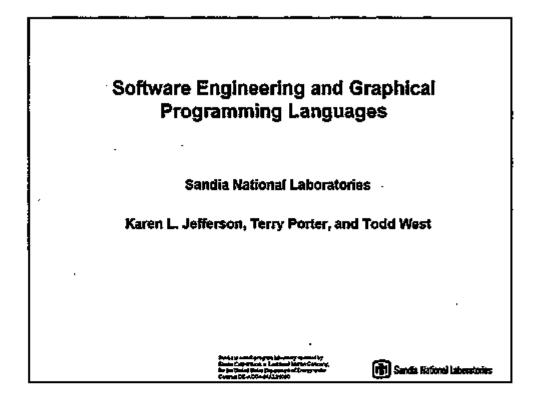
Savannah River Sile / OPS-DTB-97-0001

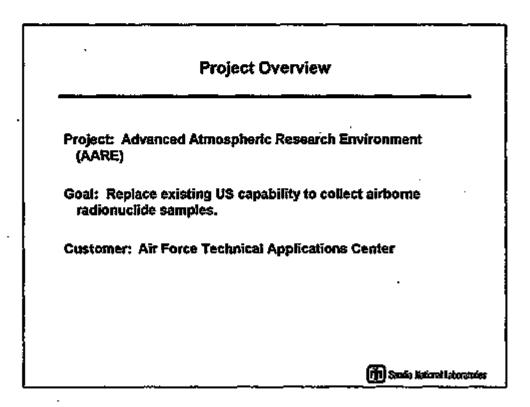
20

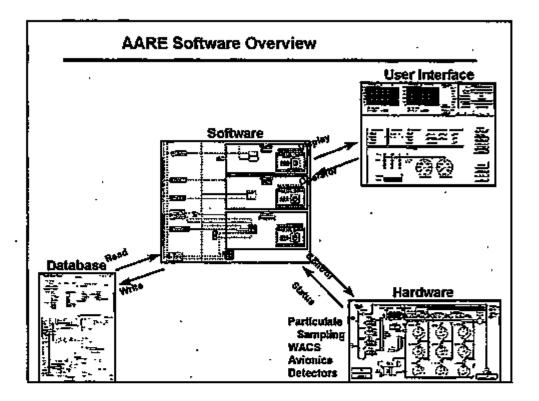
19

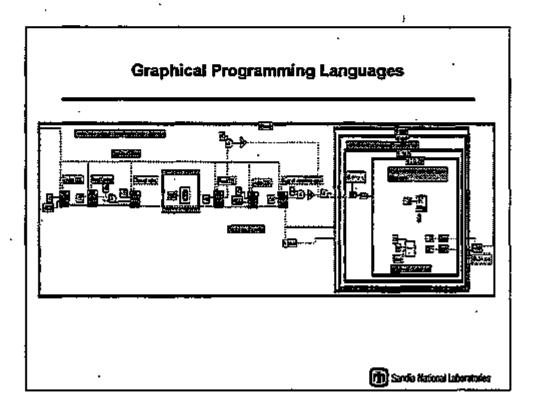


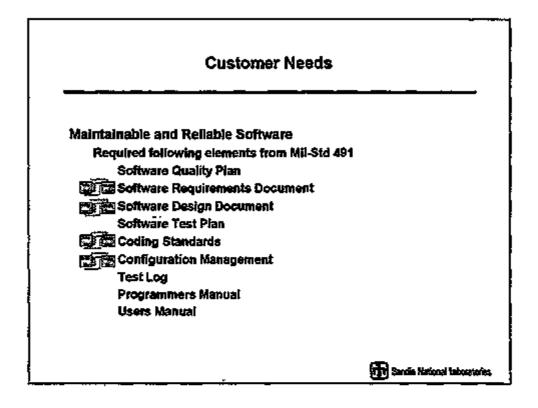
Sevenash River Scie / OPS-DTB-97-0001

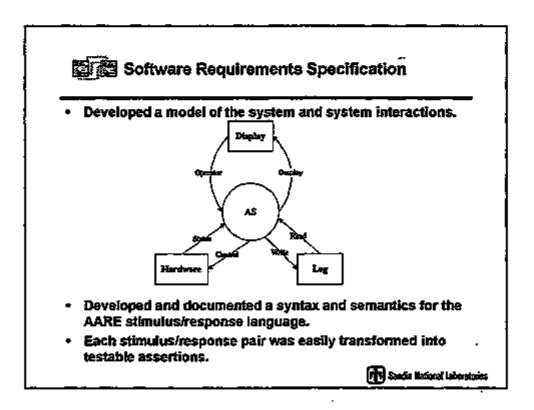


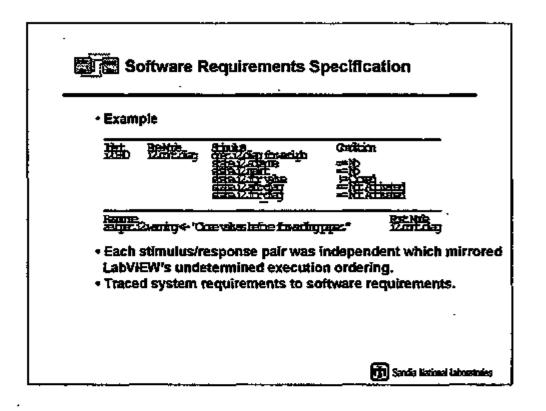


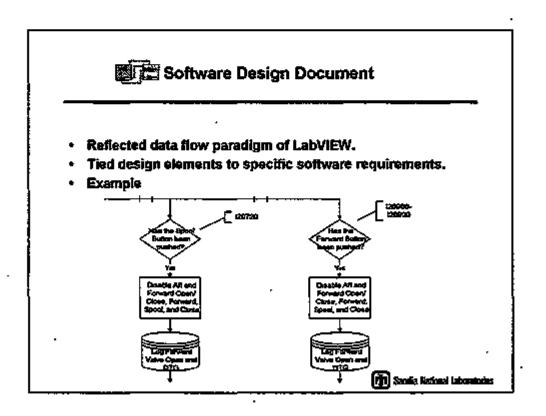




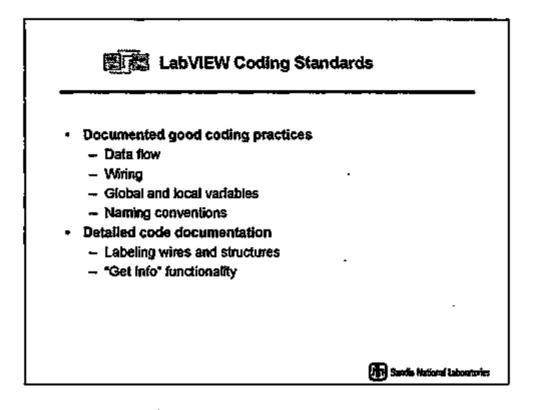


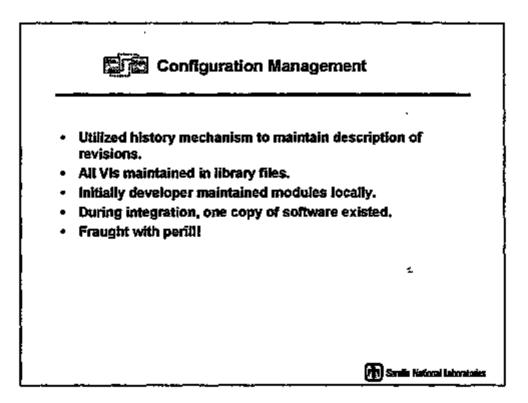


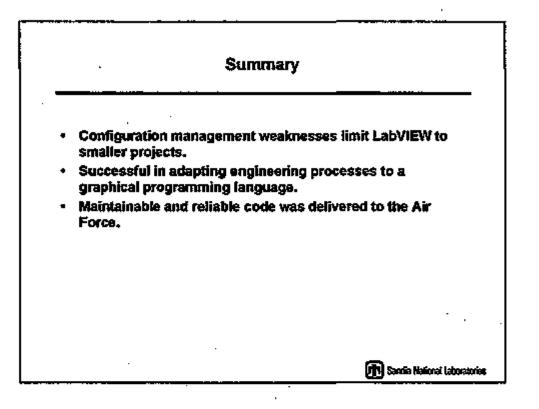




)





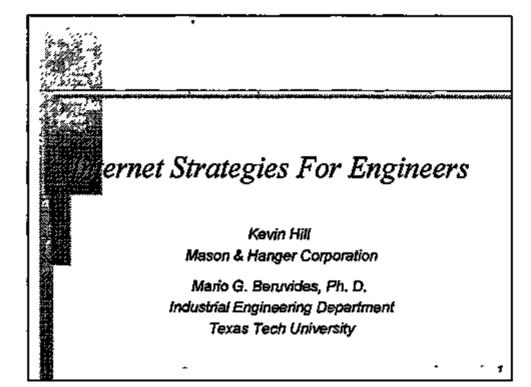


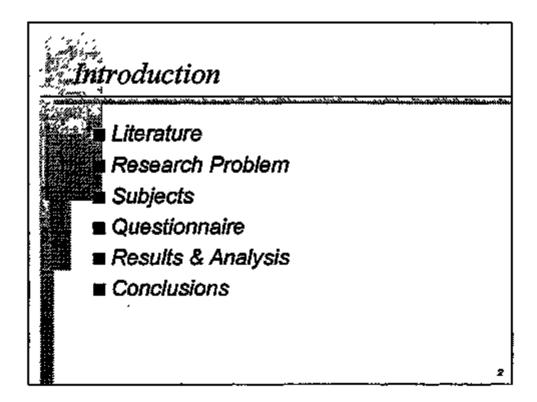


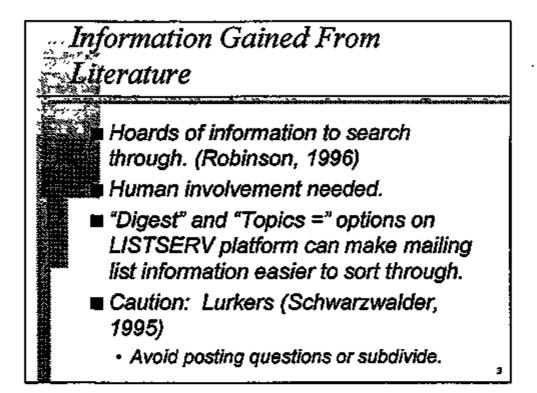
Session B2: Internet WEB Applications

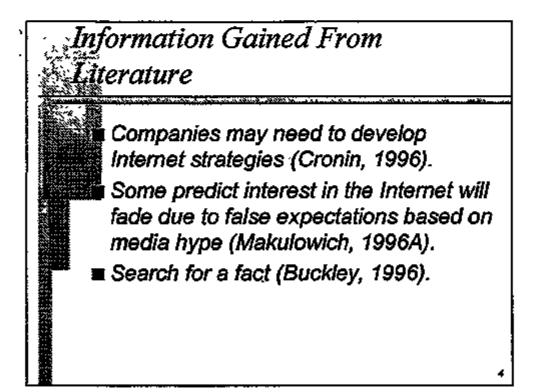
Chair Faye Brown Lockheed Martin Energy Systems

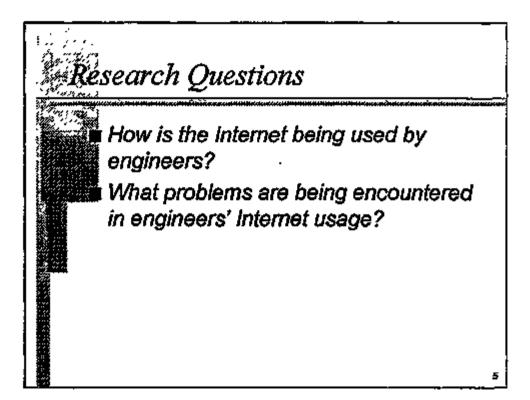
Session : Paper #	Author(s)	Title
B2:1	Kevin Hill Pantex Plant	Internet Strategies for Engineers
B2:2	David Leong & Fran Current Sandia National Laboratories	Exploiting the Intranet: A New Architecture for Enterprise Information
B2:3	Jennie Negin Sandia National Laboratories	"Rightsizing" Software Quality for a Web Services Organization

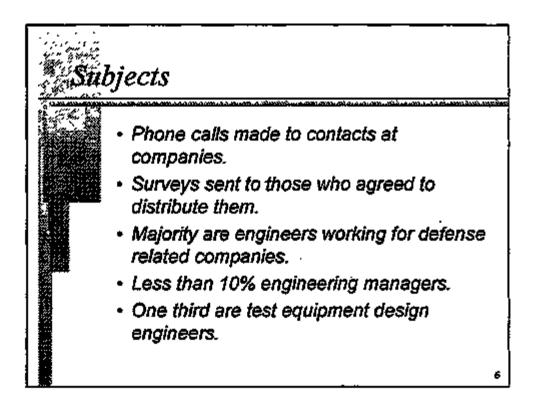


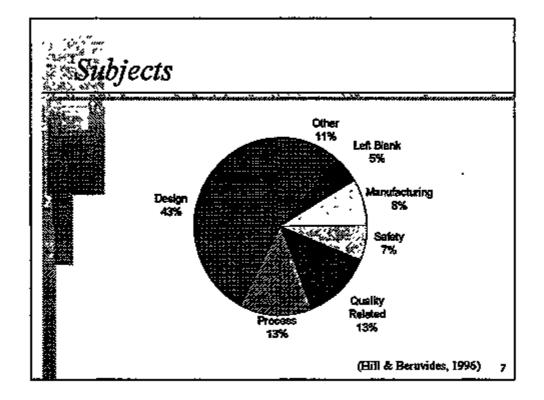




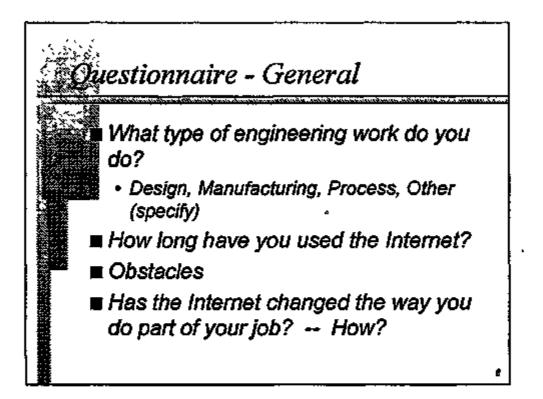


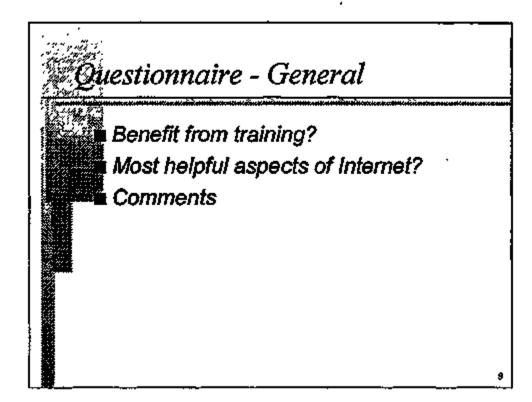




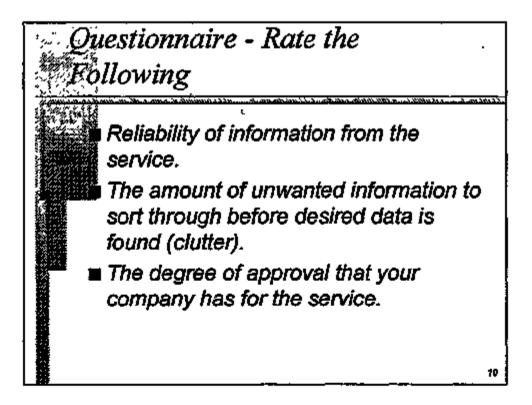


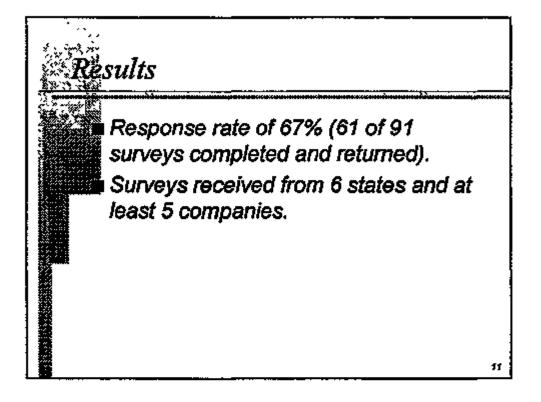
•)

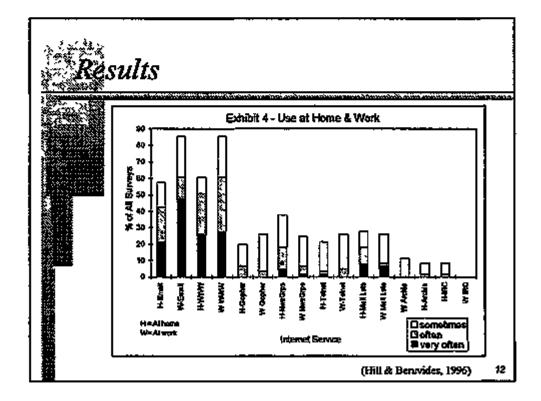




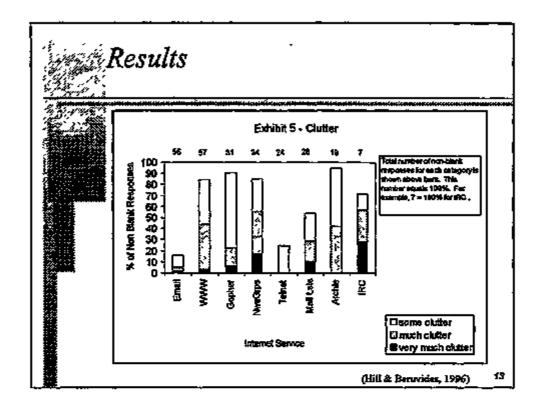
-5

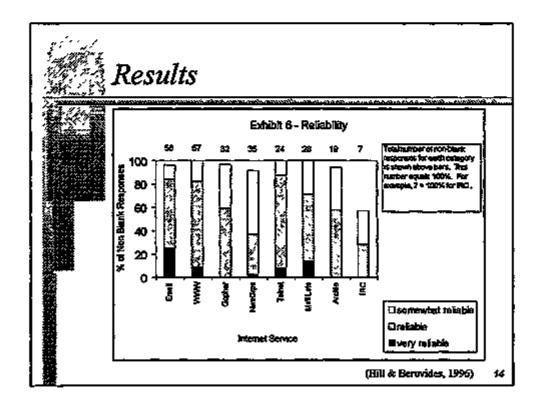




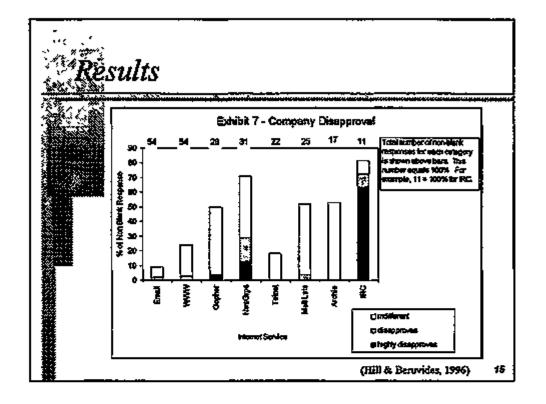


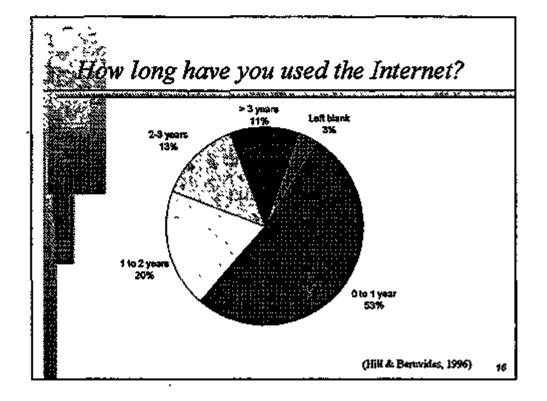
-- ,





.



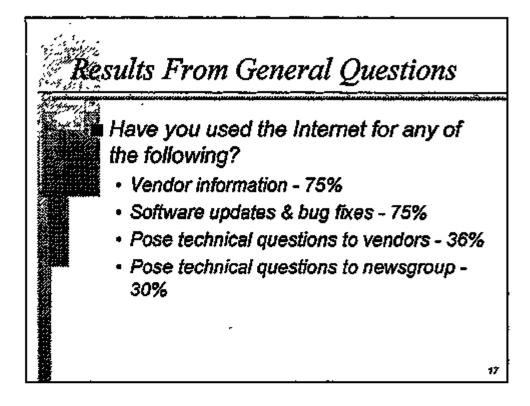


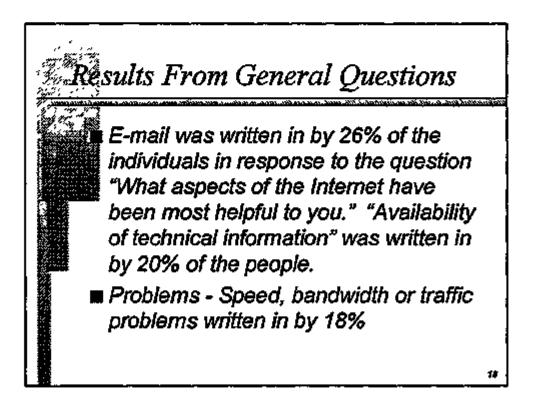
8

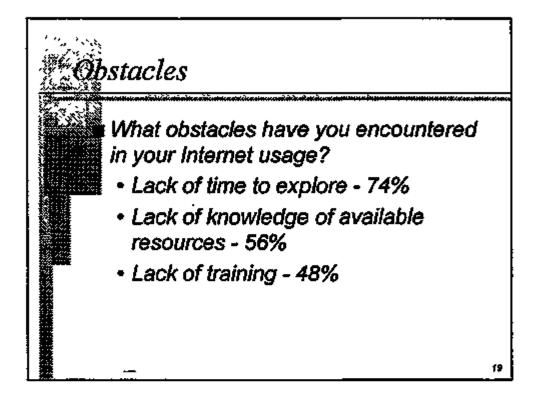
)

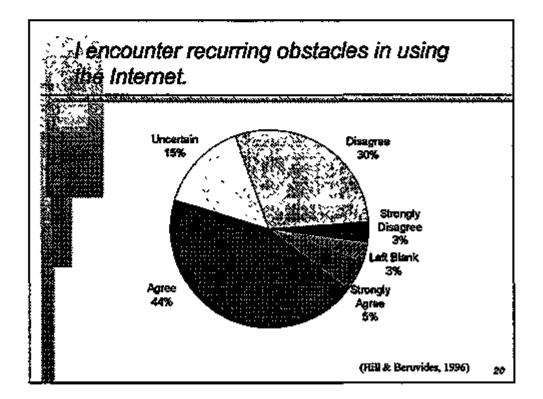
Ĵ

L.

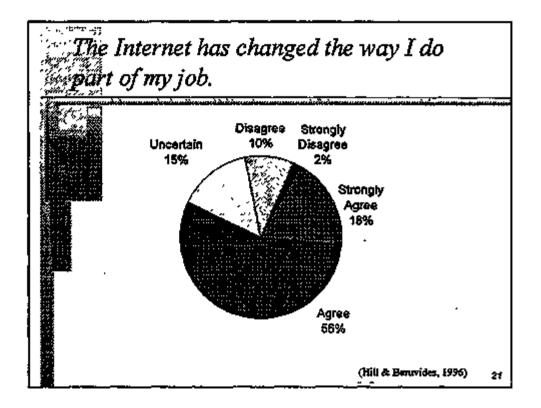


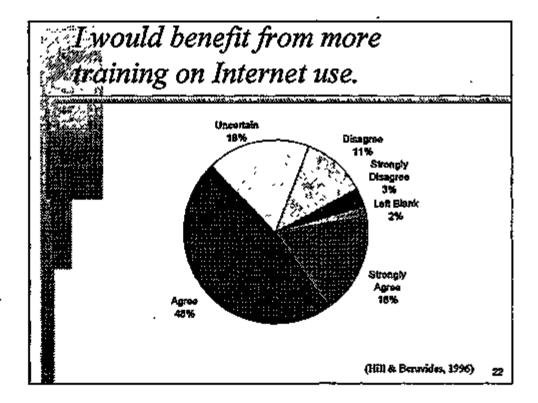


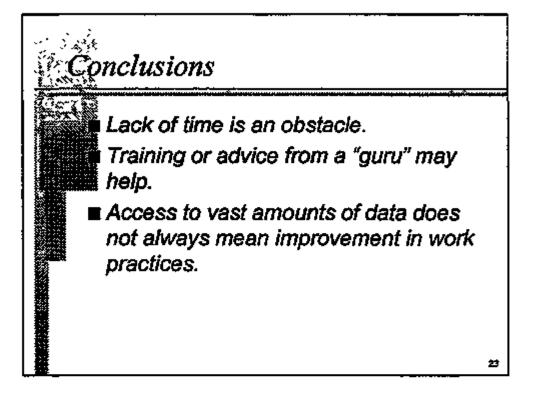


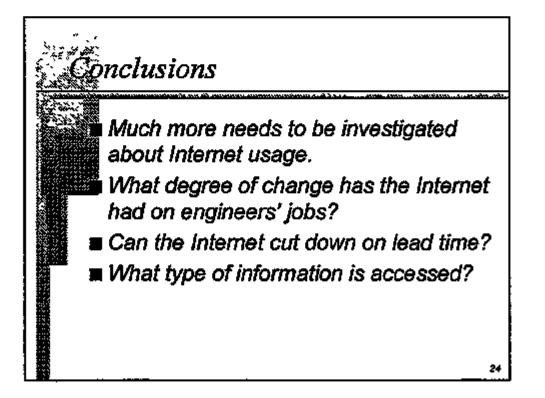


)



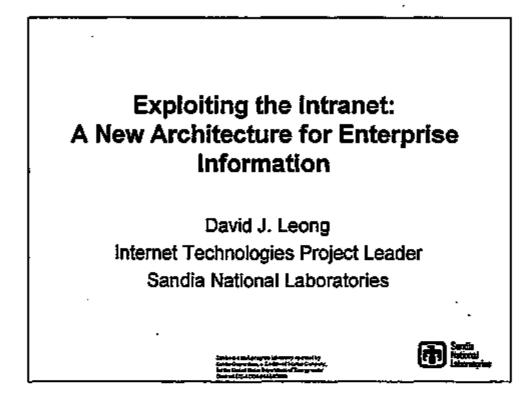


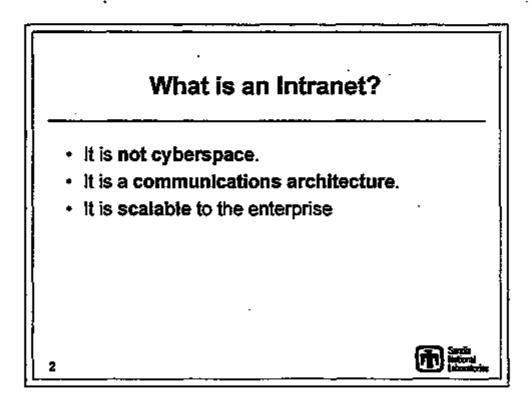


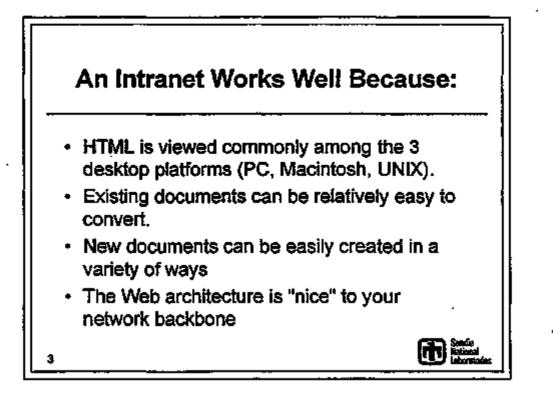


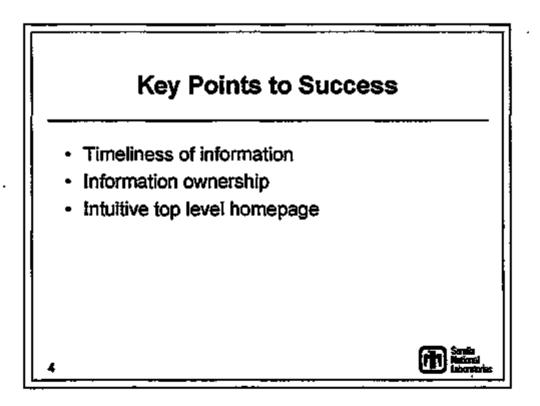
REFERENCES

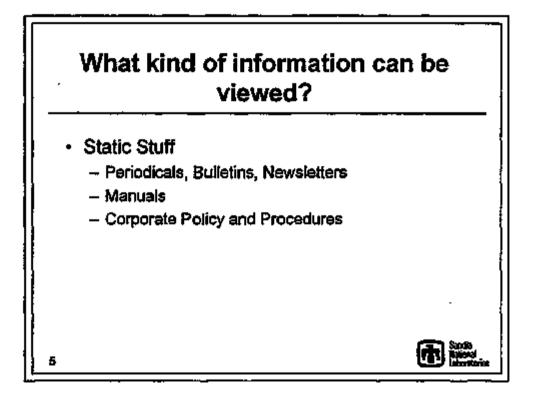
- Buckley, W. F. (1996). Is the Internet really filled with endless wonders? <u>Amarillo Daily News</u>, May 6, p. 4A.
- Cronin, M. J. (1996). <u>Global Advantage on the Internet: From</u> <u>Corporate Connectivity to International Competitiveness</u>, New York: Van Nostrand Reinhold.
- Makulowich, J. S. (1996A). Net sitings: Future trends on the net. <u>Online</u>, January/February, pp. 37-38.
- Robinson, K. L. (1996). People talking to people: Making the most of Internet discussion groups. <u>Online</u>, January/February, pp. 27-32.
- Scinwatzwalder, R (1995). Engineering and the Internet: A survivor's Manual. <u>Database</u>, April/May, pp. 72-74.
- Hill, Kevin and Bernvides, Mario, "Strategies For Coping With The Internet: A Survey of Engineers' Usage And Problems" <u>Proceedings of the 1996 National Conference of the American Society for Engineering</u> <u>Management</u>, pp. 317-323

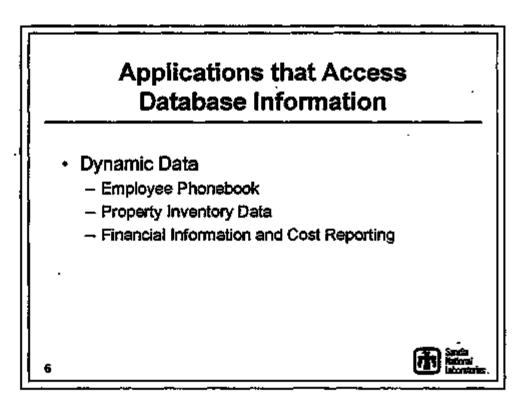


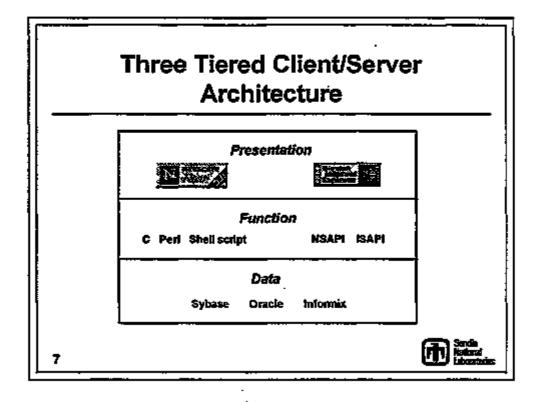








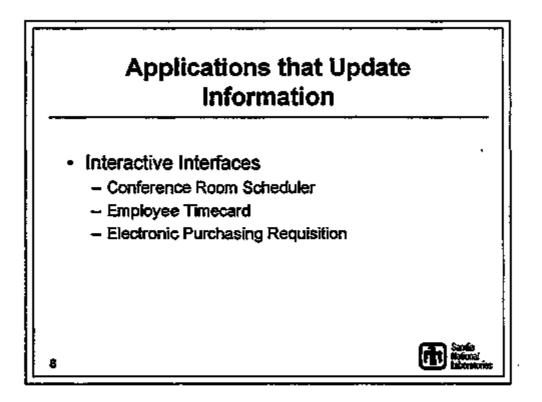


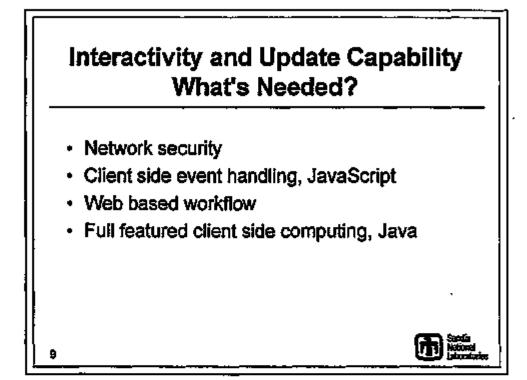


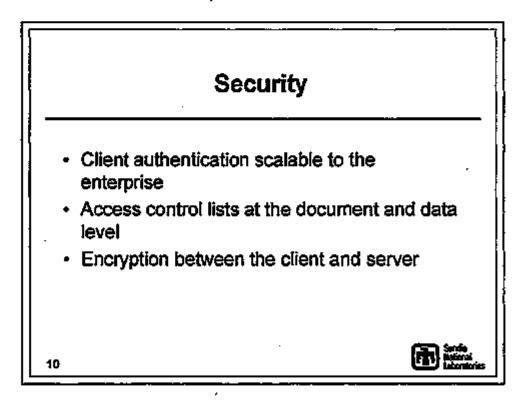
ົ)

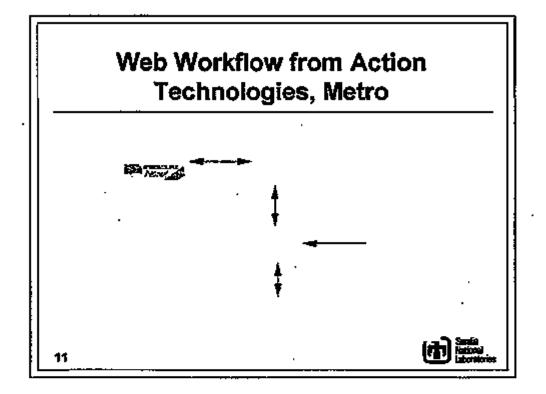
)

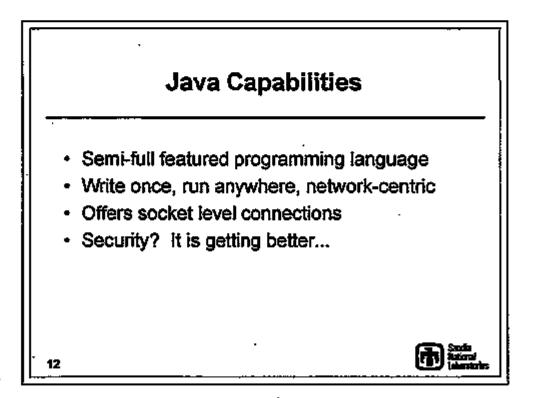
·)

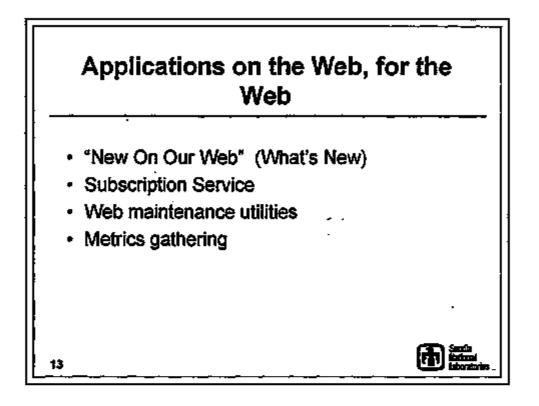


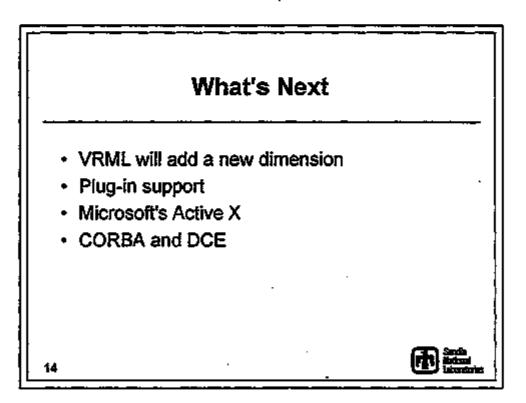


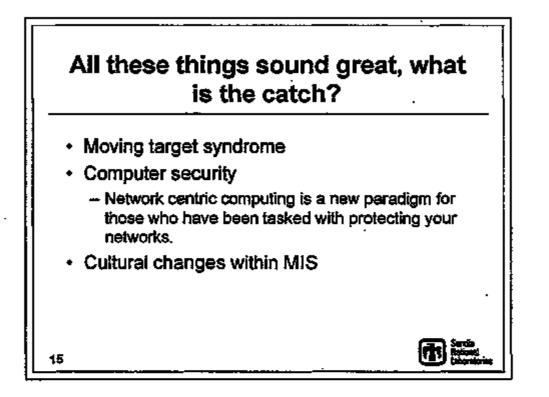


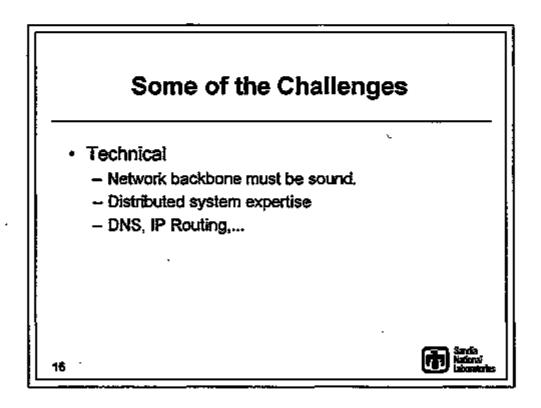


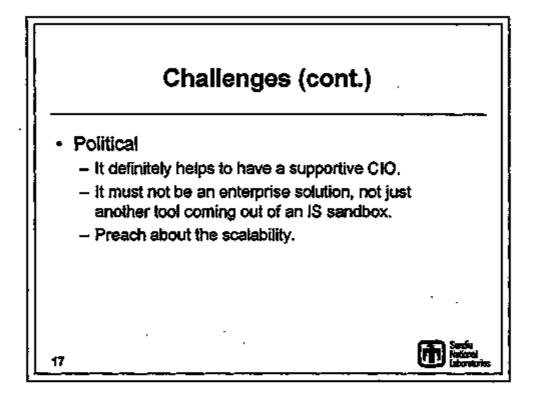


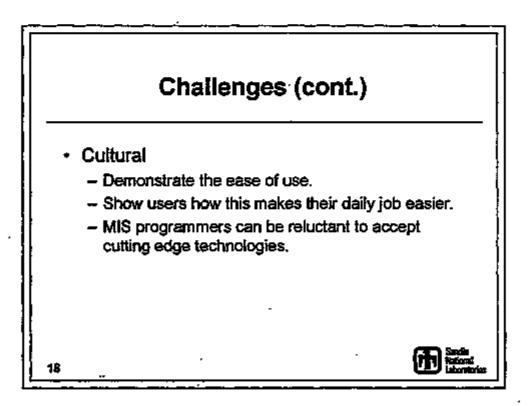


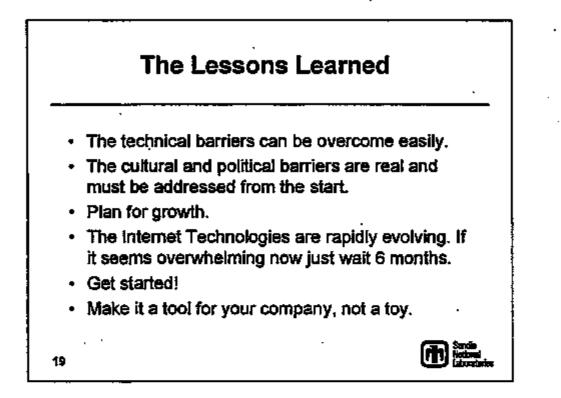


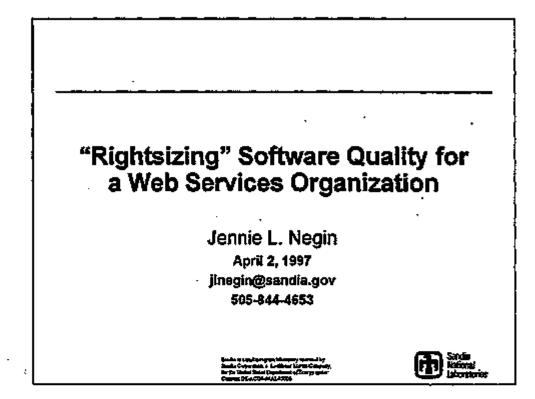


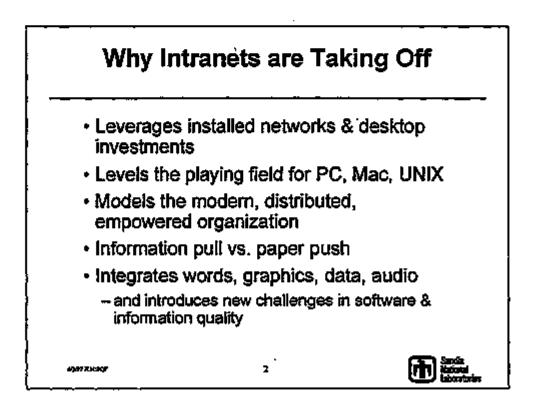


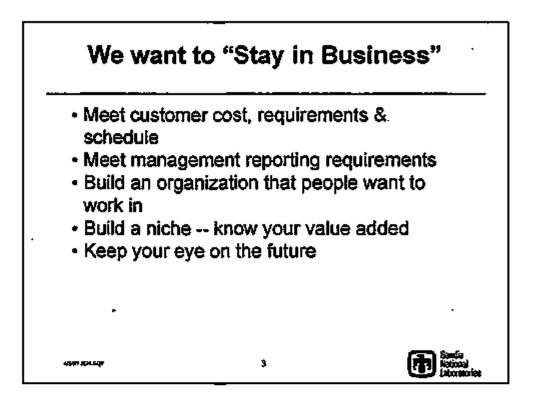


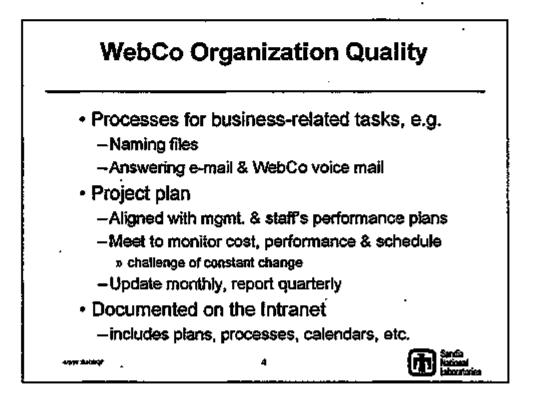




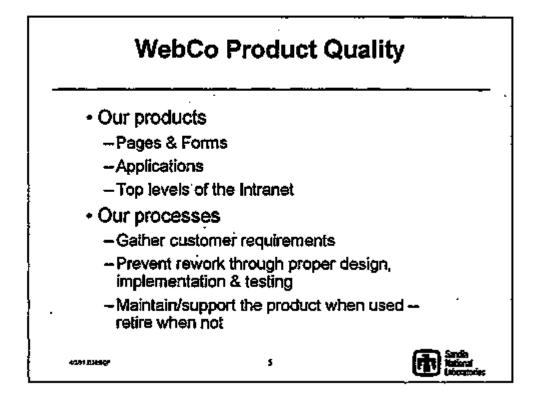


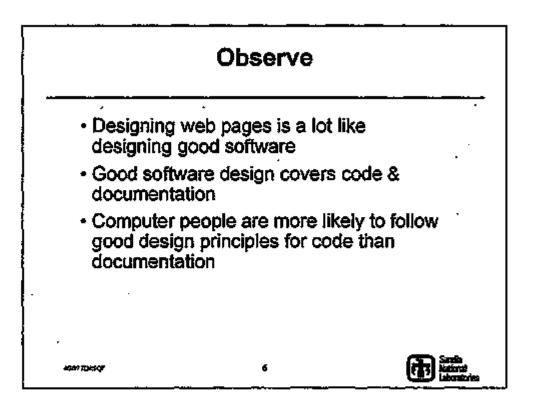


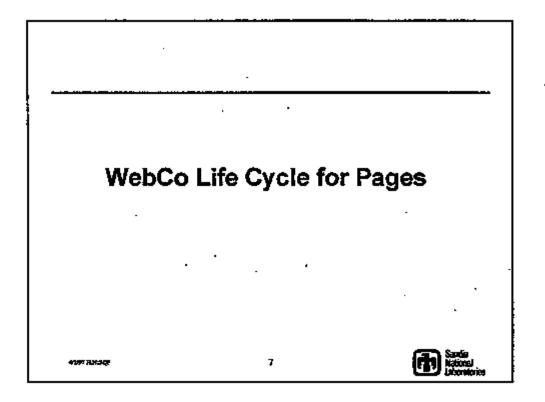


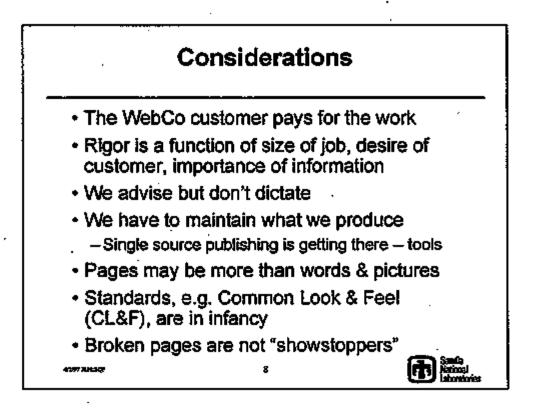


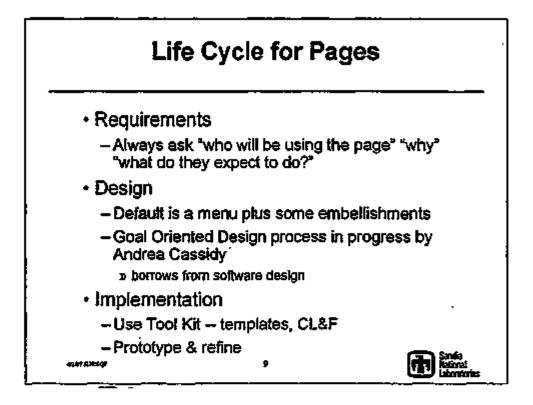
--)

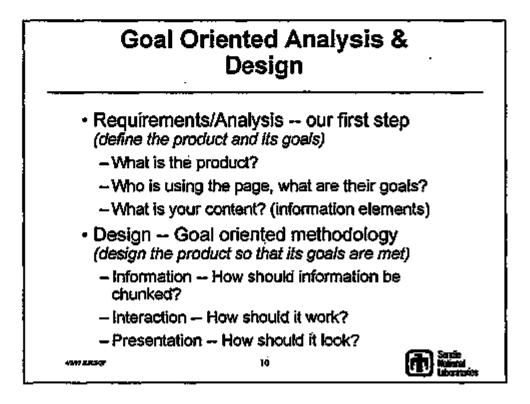


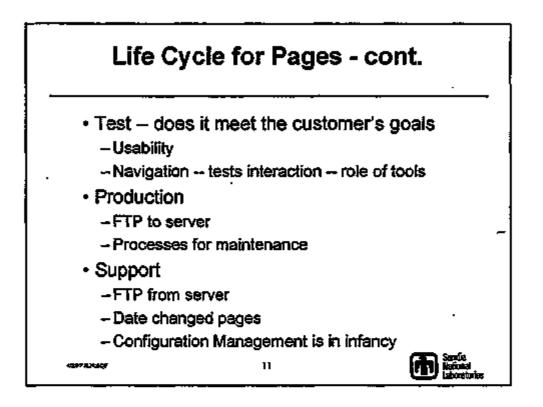


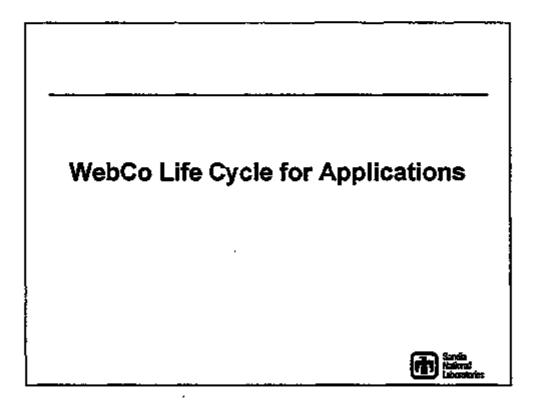




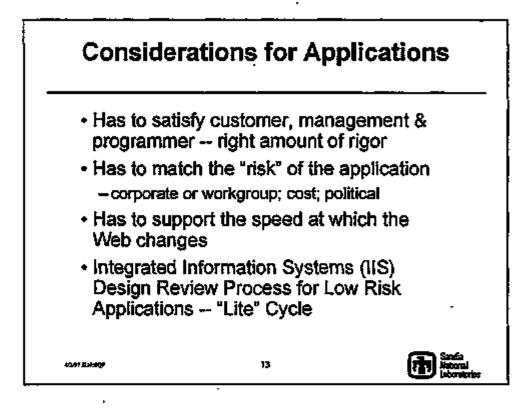


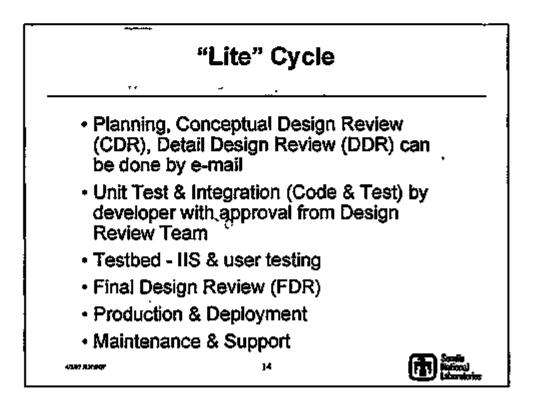


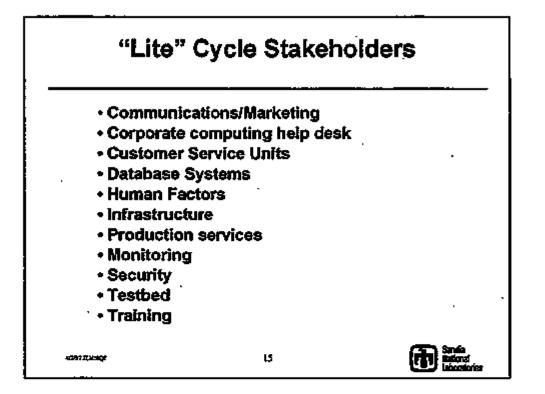


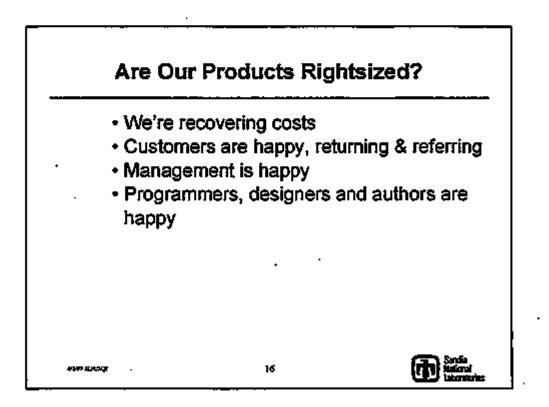


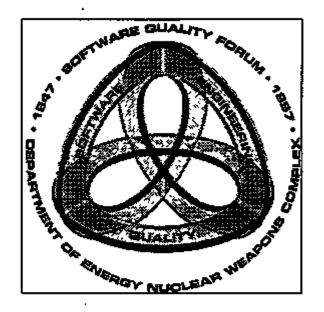
)









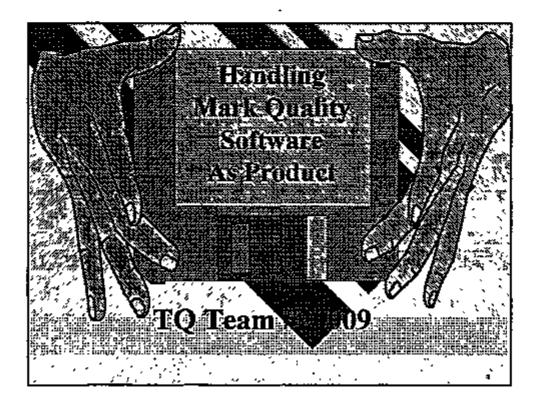


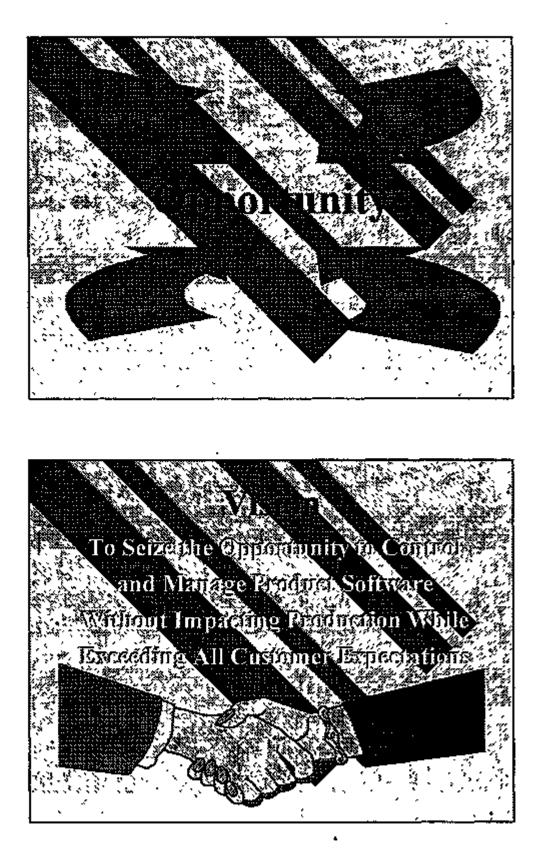
Session A3: Software Process Improvement I

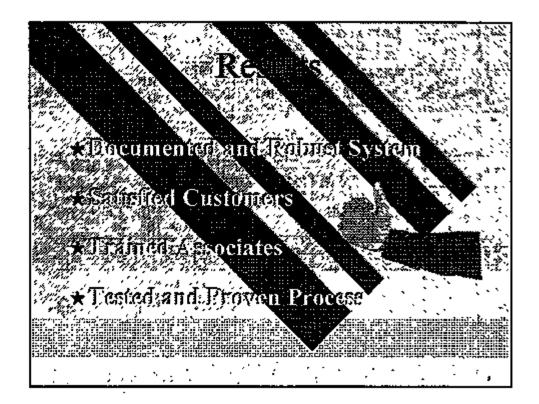
Chair Mike Lackner AS/FM&T

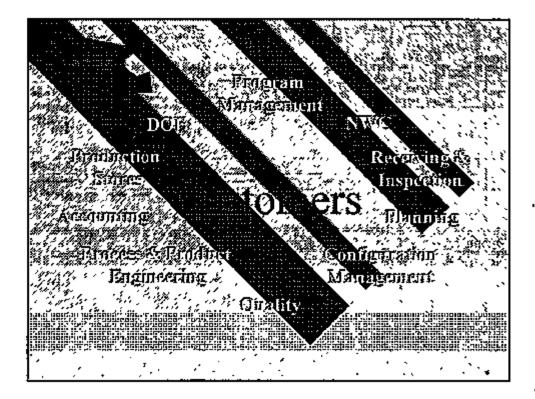
Session : Paper #	Author(s)	Title
A3:1	Don Schilling	Quest for Excellence 1996: Reaching
	AS/FM&T	for the Stars
A3:2	Don Rathbun	Command Media System at the
	AS/FM&T	Kansas City Plant (KCP)
A3:3	Michael Tiemann	Departmental Information
	Headquarters Department of Energy	Architecture

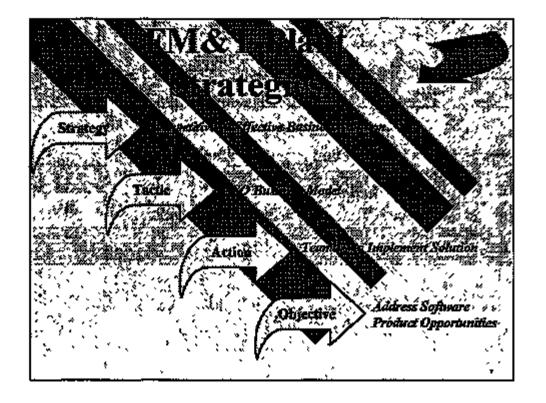


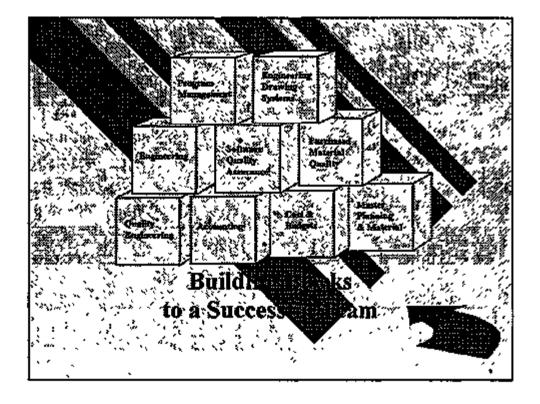








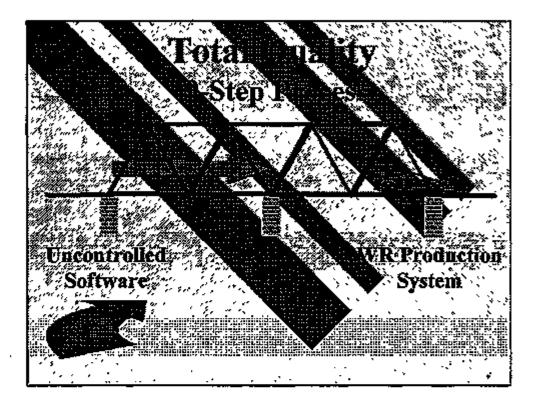


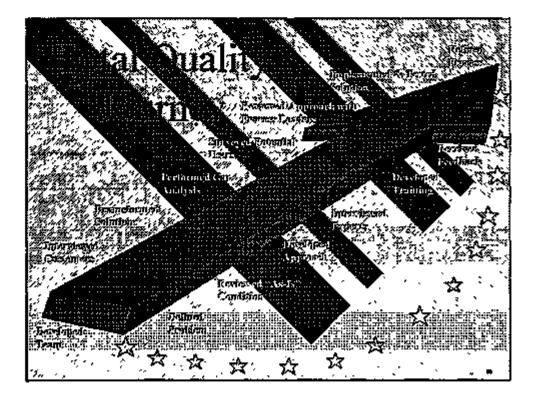


-

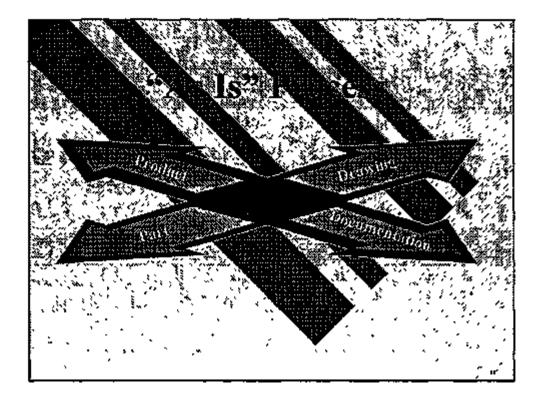
~)

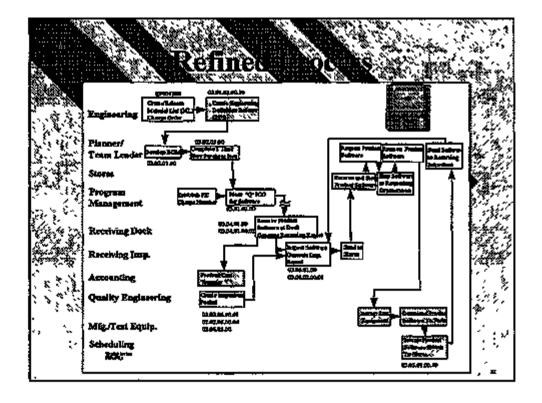
· ·)





SQAS Forizn April 1-3, 1997





6

<u></u>,

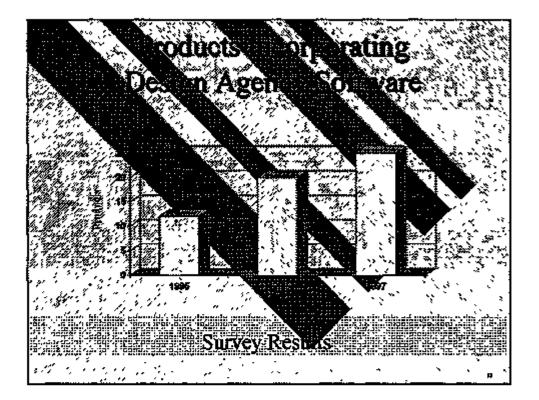
ı

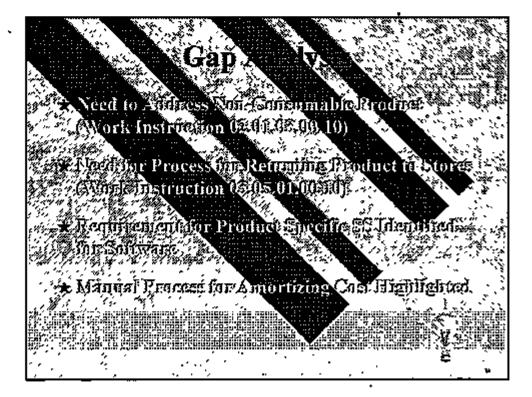
-,

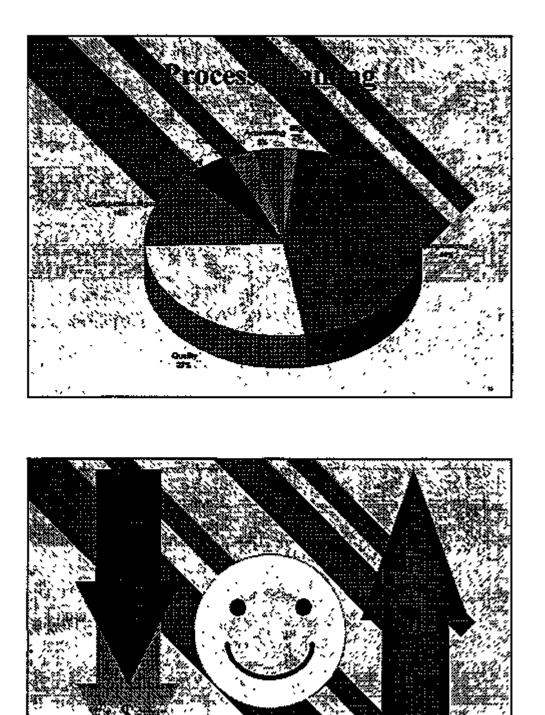
⁻)

١.

ı







Tistomer

ŝ.

SQAS Forum April 3-3, 1997

8

1

1

I

4

Õ

,

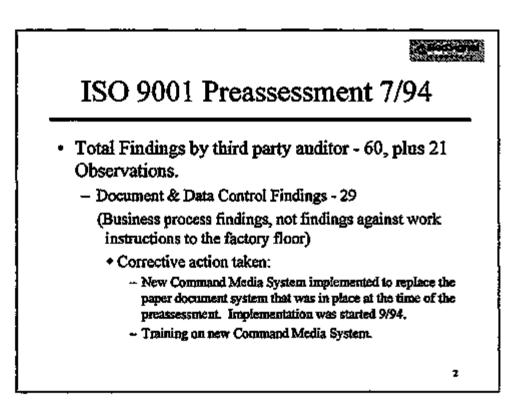
ı.

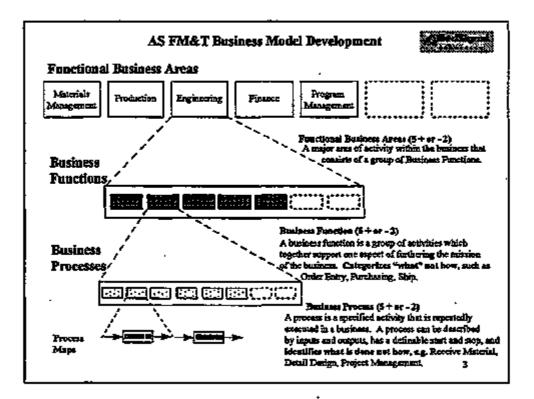
Command Media System at the Kansas City Plant (KCP)

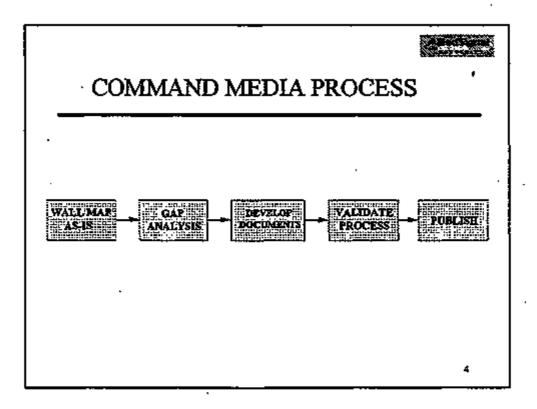
Don A. Rathbun, Staff Engineer AlliedSignal Federal Manufacturing & Technologies (FM&T)*

Presented at the 1997 Software Quality Forum, April 1-3, 1997 Kintland Air Force Base, Albuquerque, New Mexico Sponsored by Department of Energy (DOE) Quality Managers Software Quality Assurance Subcommittee of the DOE Quality Managers Weapons Quality Division, DOE-Albuquerque Office

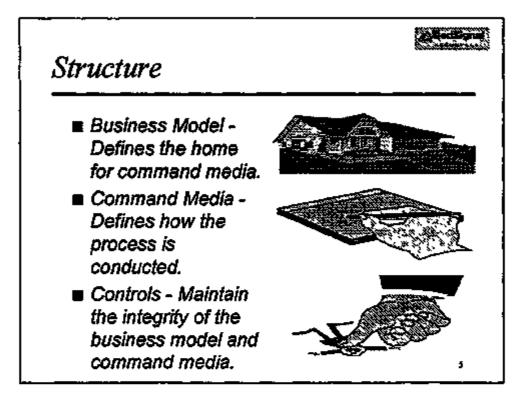
*Operated for United States Department of Energy under Contract No. DE-AC04-76-DP00613 ©Copyright AlliedSignal Inc., 1997.

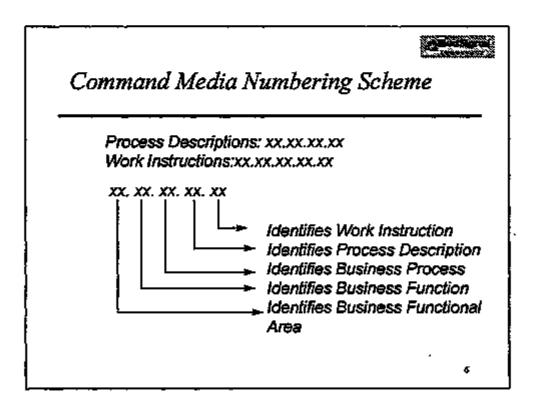


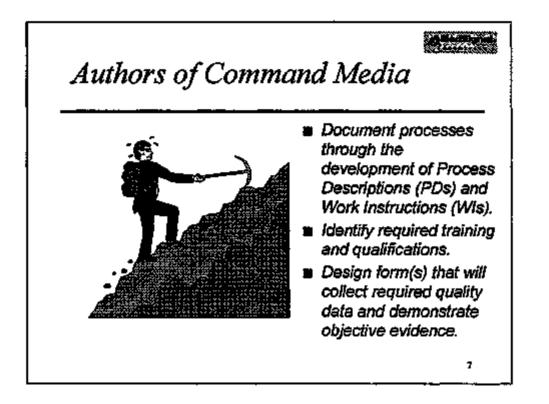


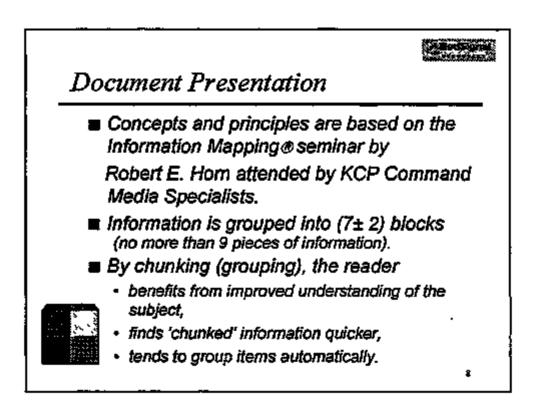


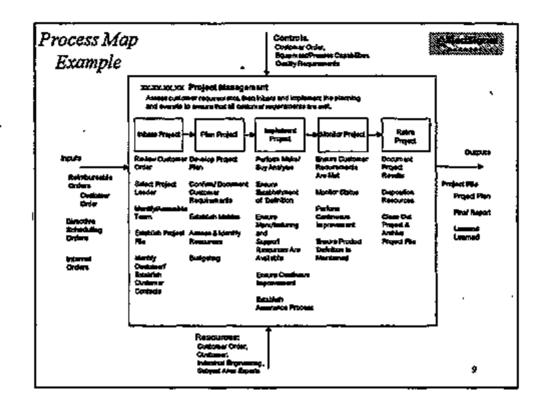
ſ

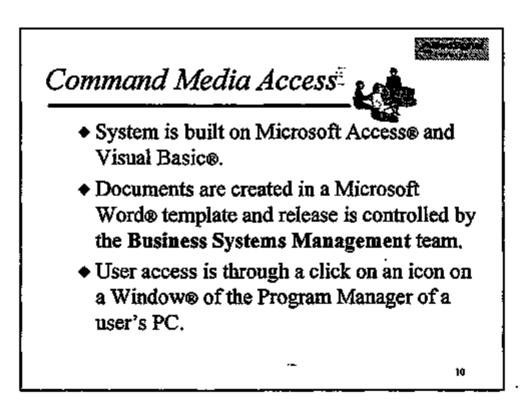


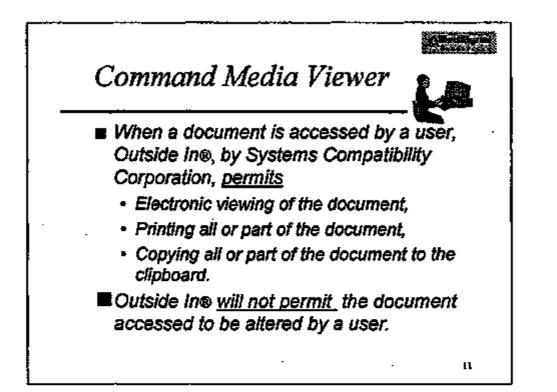




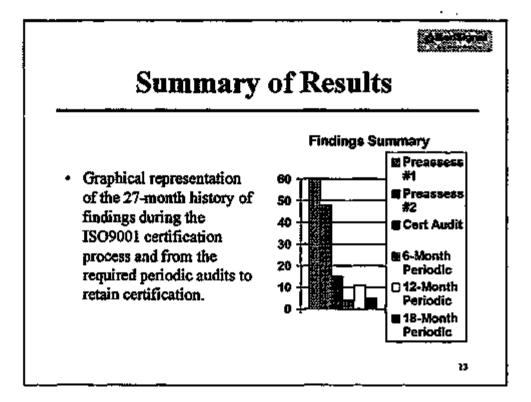


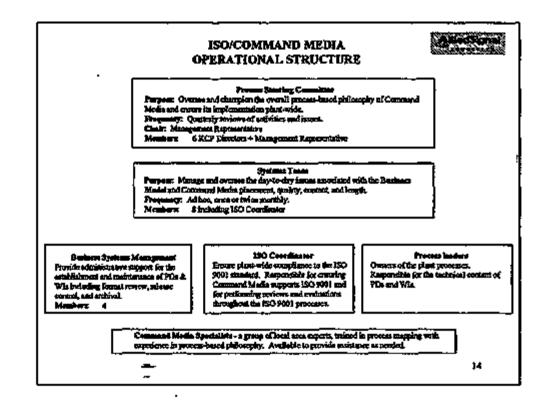


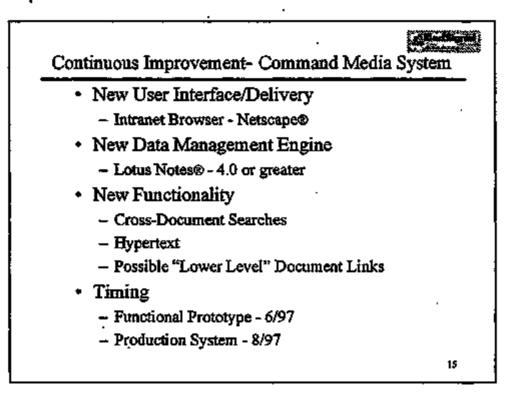


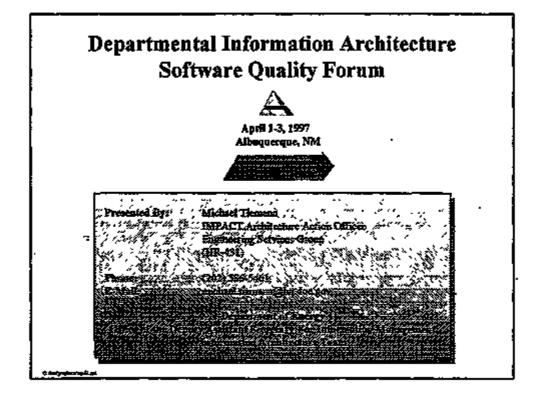


Summary of Results After Implementing The New Command Media		
Assessment	<u>Total Findings</u>	Document & Data Control Findings
Preassessment 7/94	60	29
Preassessment 2/95	48	23
Certification 4/95	15	3
6-Month Periodic 10/95	4	1
12-Month Periodic 5/96	11	3
18-Month Periodic	5	1
10/20		12



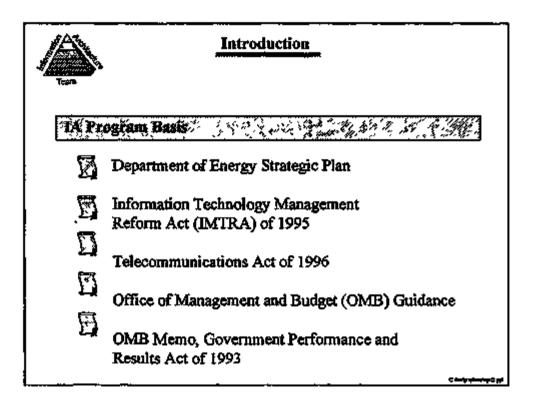


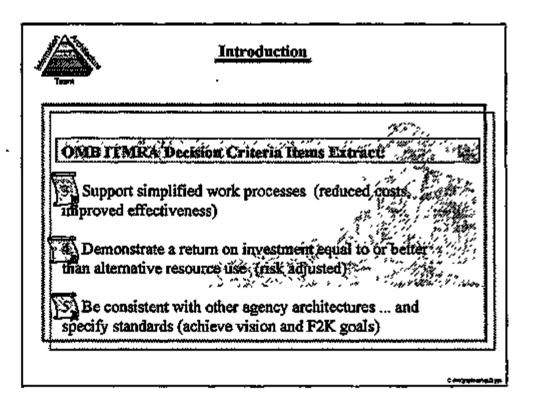


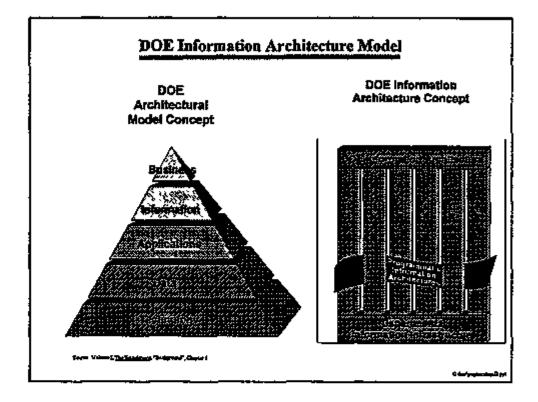


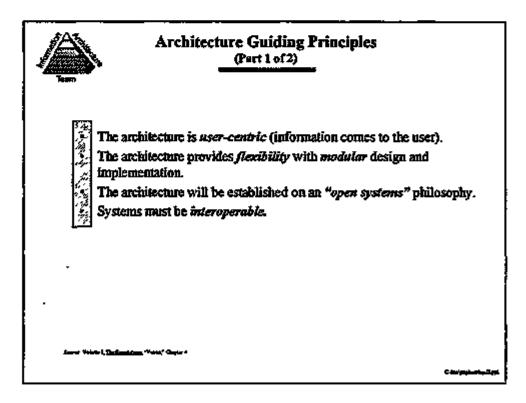
<u>Outline</u> Information Architecture Program Introduction Models and Principles - Publications Future Directions DOE IA Guidance Highlights Software Implications Discussion

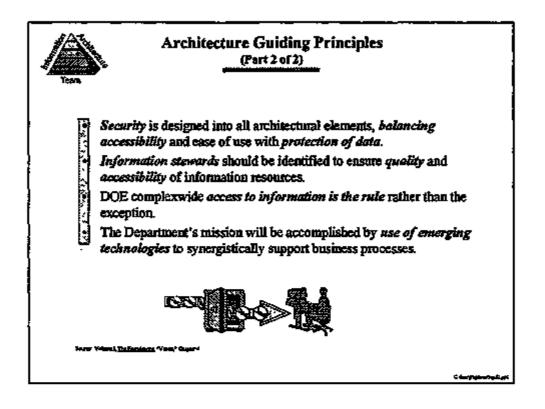
I

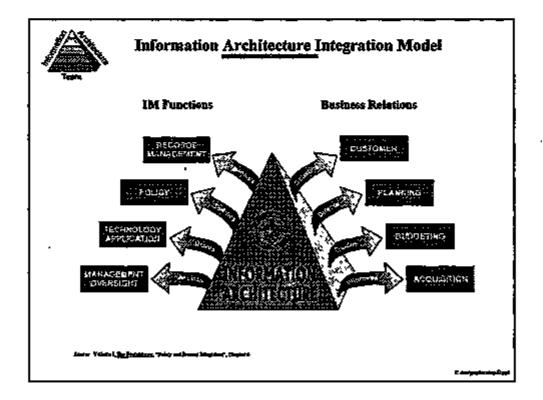


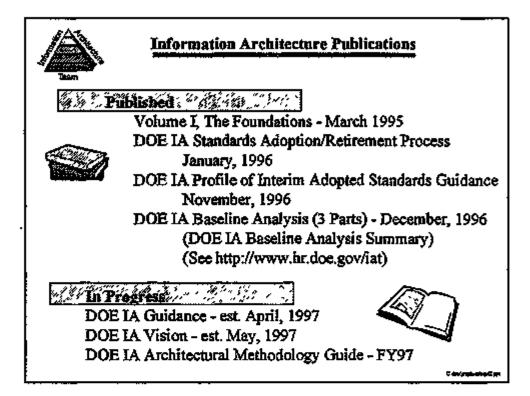


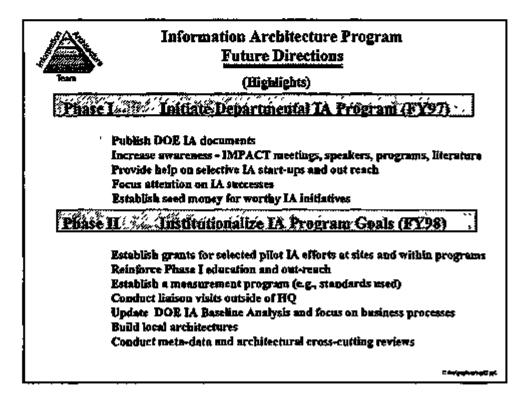


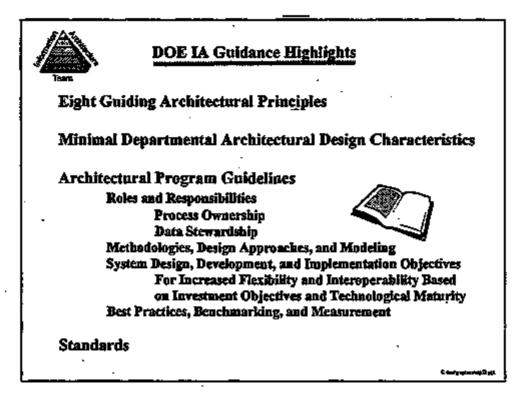


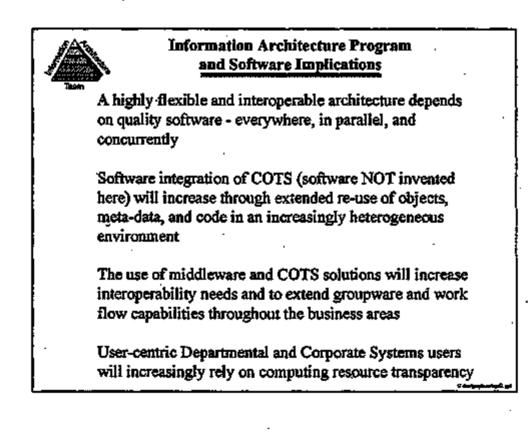




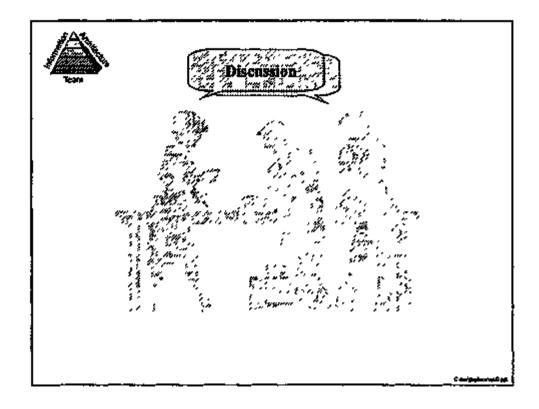








б



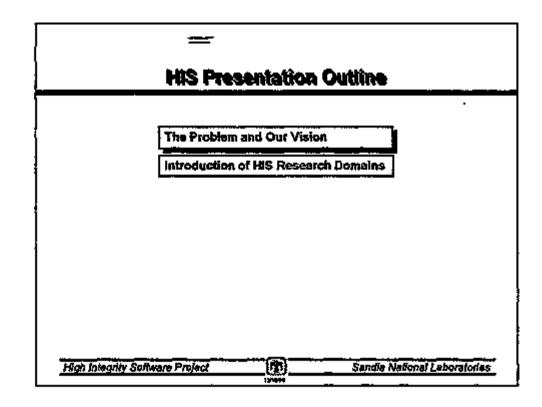


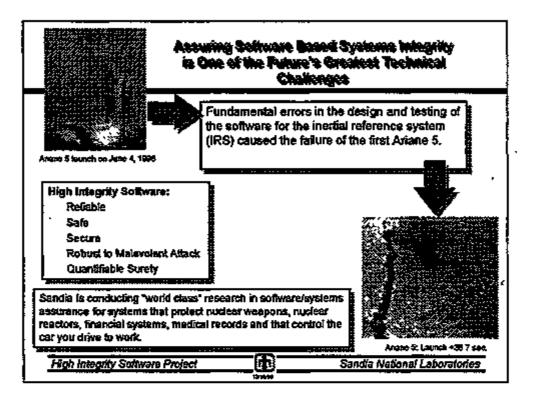
Session B3: High Integrity / Formal Methods I

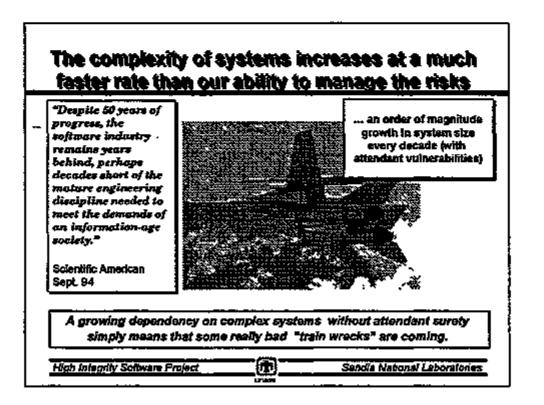
Chair Dave Peercy Sandia National Laboratories

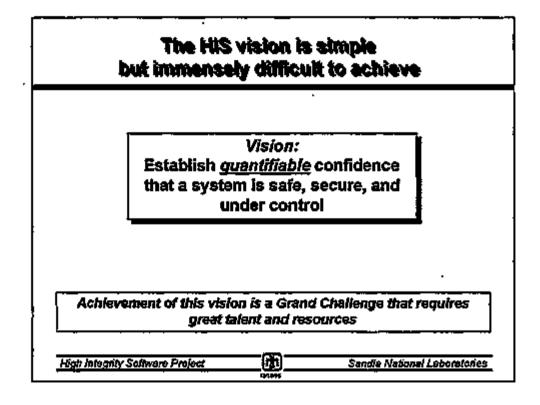
Session : Author(s) Title Paper # Larry J. Dalton & Marie-Elena Kidd Meeting the High Integrity Software B3:1 Sandia National Laboratories Needs of Today and Tomorrow Victor Winter An Overview of the AST Software B3:2 Sandia National Laboratories Construction Methodology Alex Yakhnis & Vladimir Yakhnis Towards Automated Construction of B3:3 Dependable Software/Hardware **Pioneer Technologies** Systems

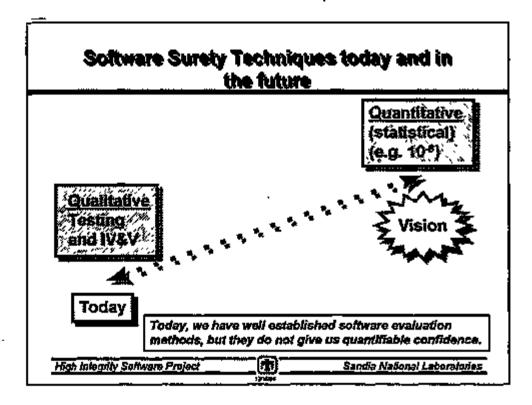


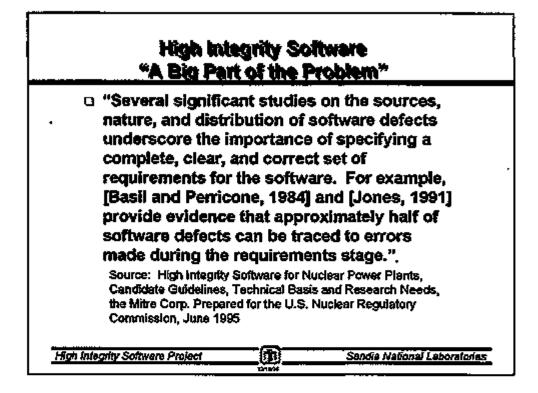






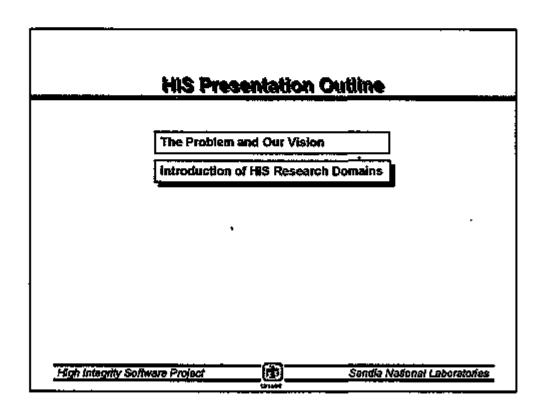


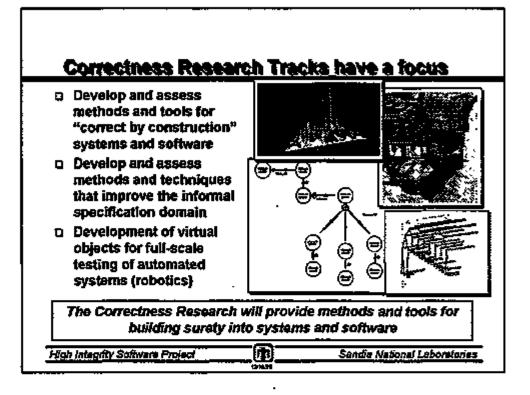


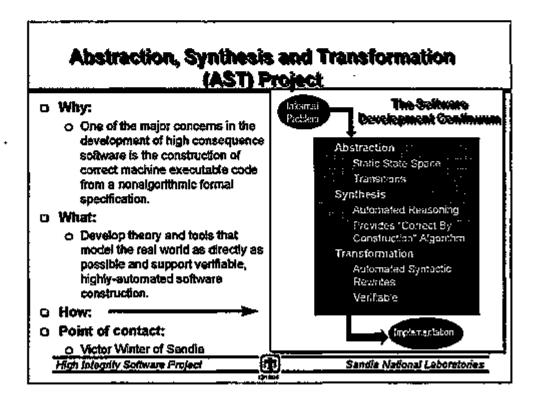


I

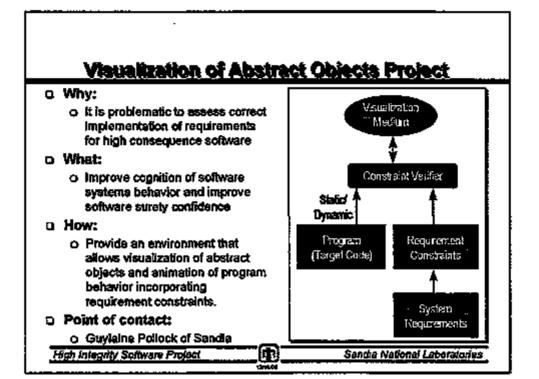
L

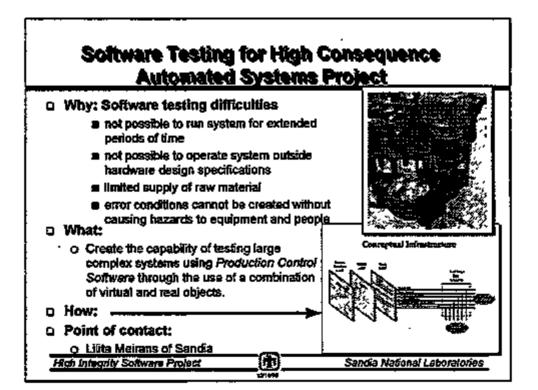






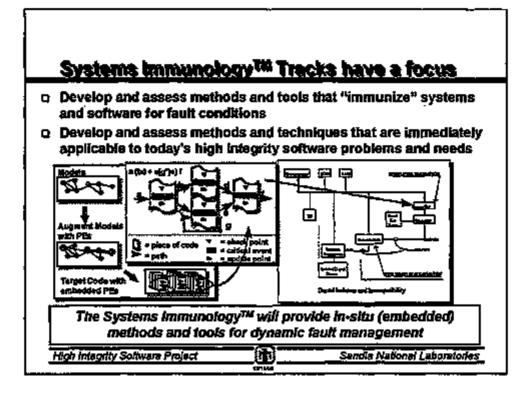
\$

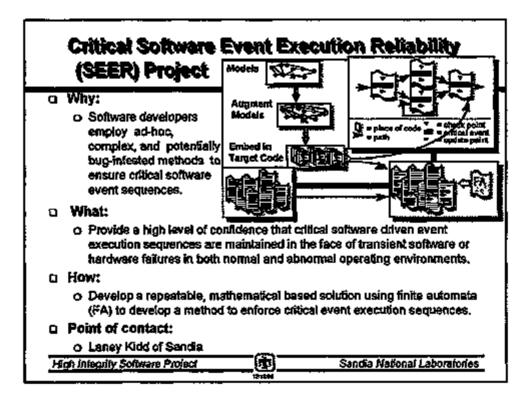


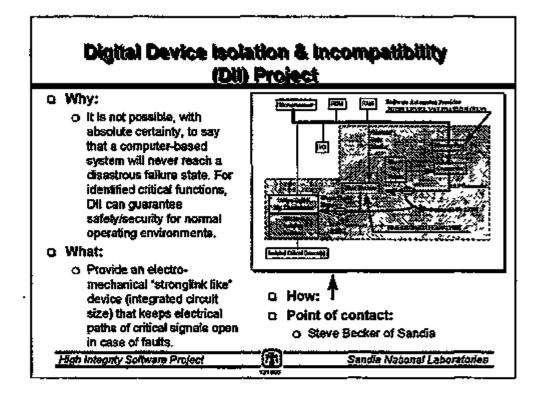


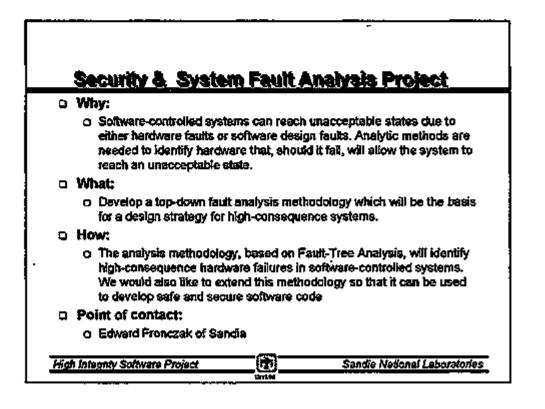
L

L





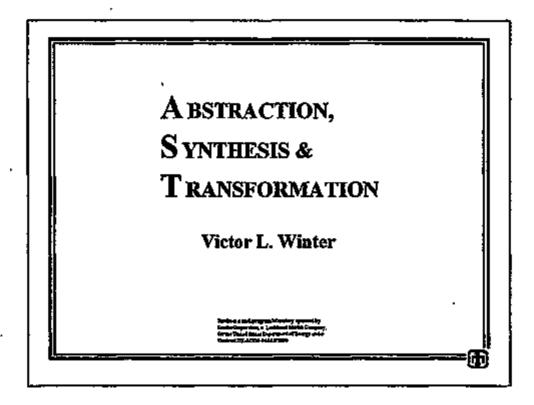


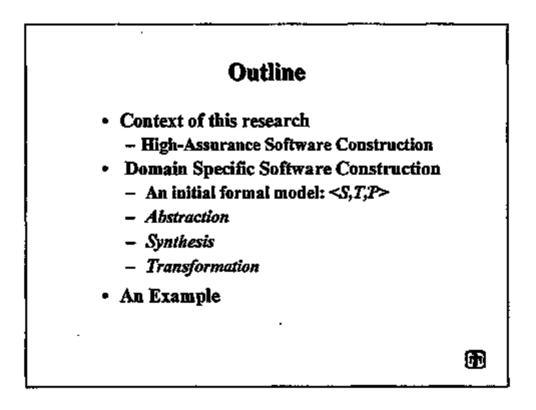


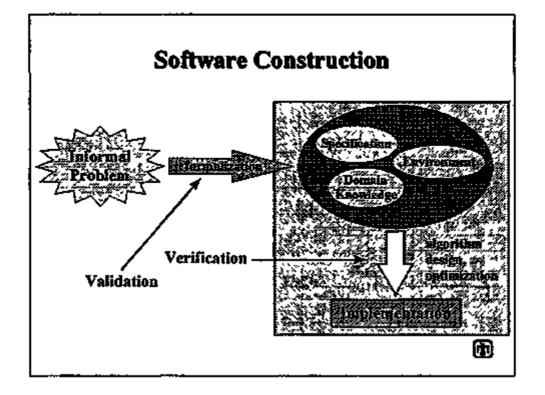
Systems with unknown surety will continue to be built based on best effort

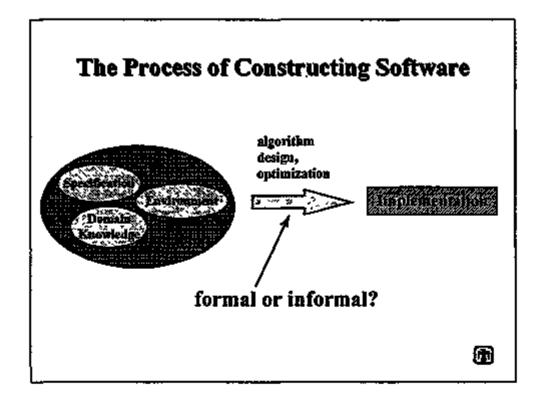
Quantifiable confidence in software-based systems is a monumental task that has been underway for many years without great success (it's really hard). New approaches, new application of mathematics; new science, fresh ideas, great cooperation without boundaries, great determination and resources will be required for several years to achieve the vision of quantifiable confidence in software based systems.

It is absolutely essential that we commit ourselves to improving the science of software-based systems







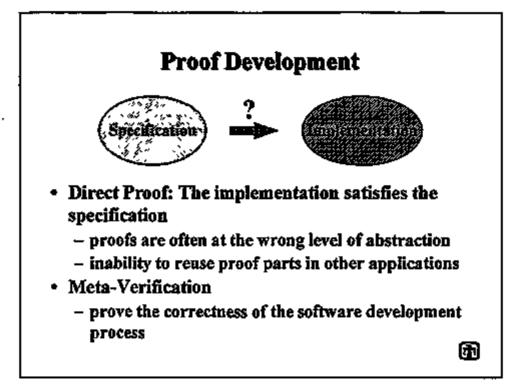


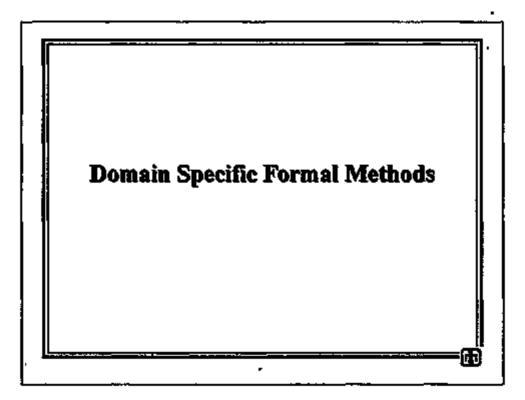
L

L.

Т

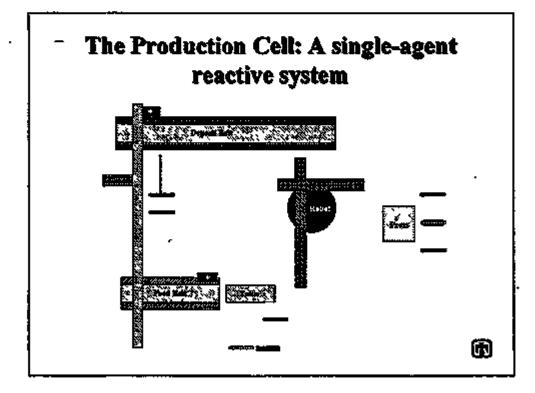
)

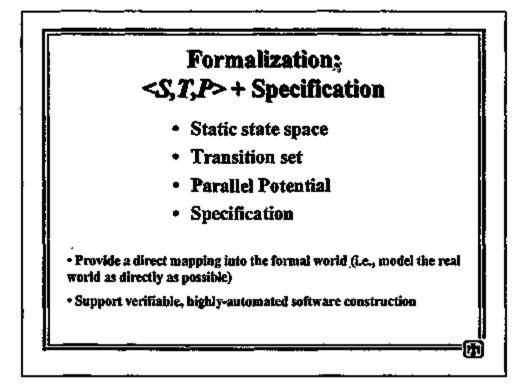


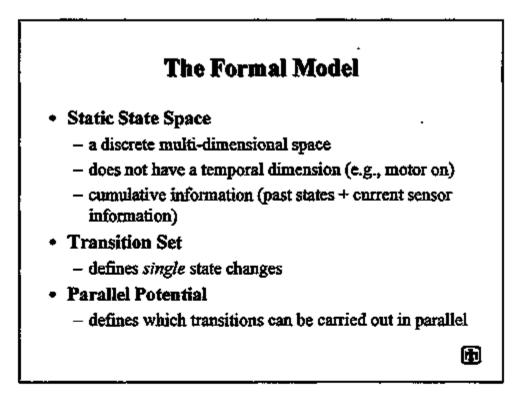


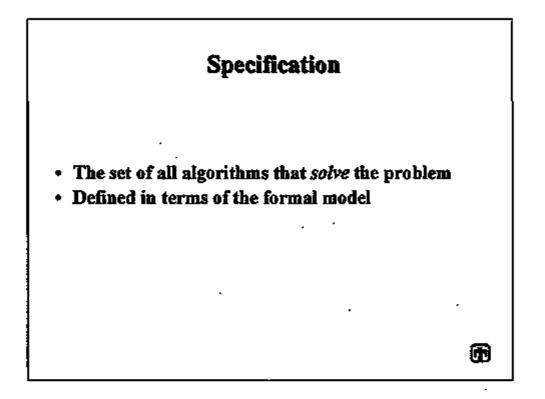
Safety-Critical Single-Agent Reactive Systems

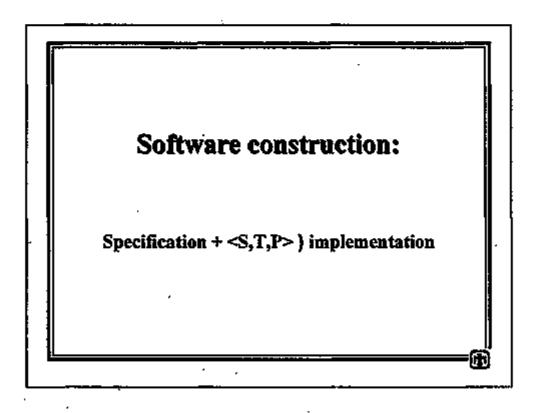
- Reactive System
 - some aspect of *time* usually plays a central role (e.g., state changes may take time—they are not instantaneous)
 - controller polls sensors to determine what state the system is in
 - parallel activities are often possible
- Single-Agent
 - all transitions initiated by the controller
 - deterministic transitions



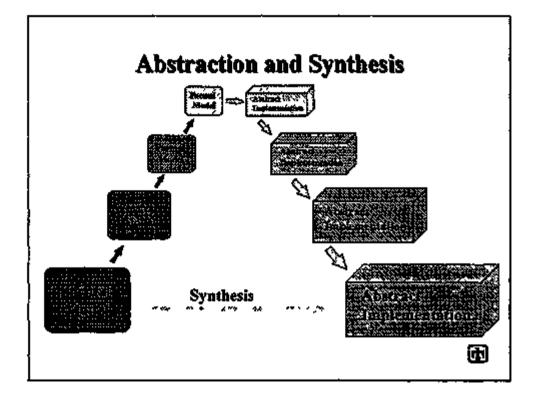


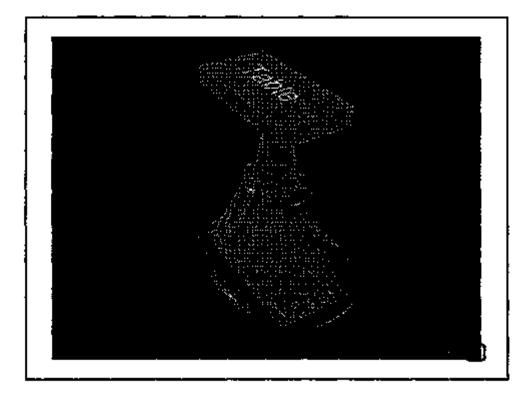


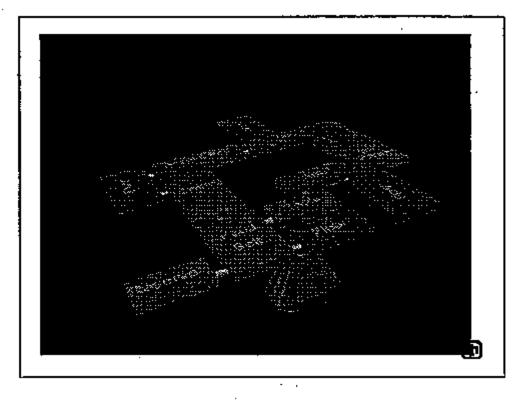


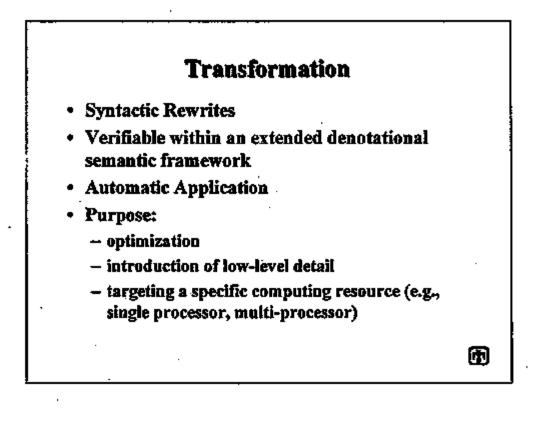


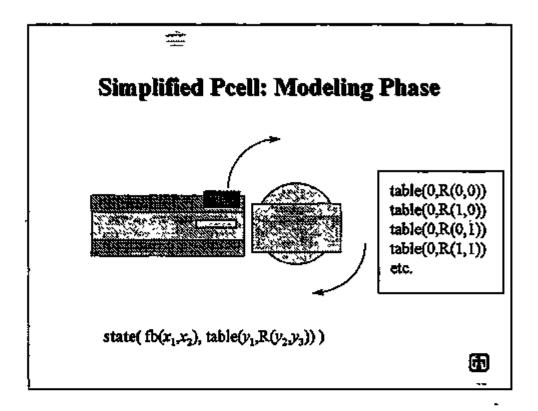
^)

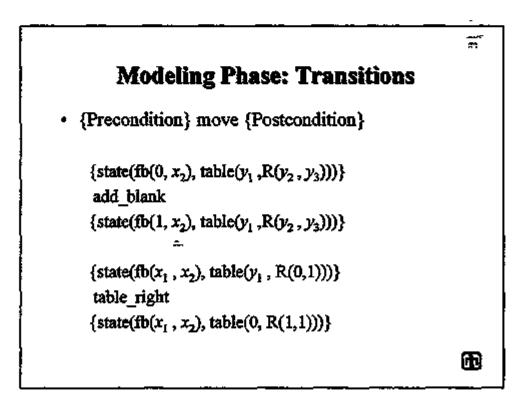














{
 (table_up, table_down),
 (table_left, table_right)
}

Construction of a Formal Specification

- Definition:
- processed a plate is processed when it "disappears" from the table
- Informal Specification:

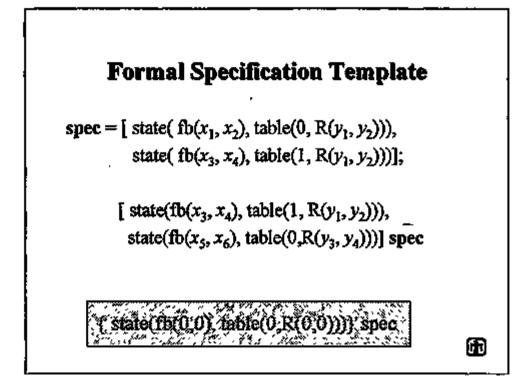
"The objective is to make the system process an infinite number of plates."

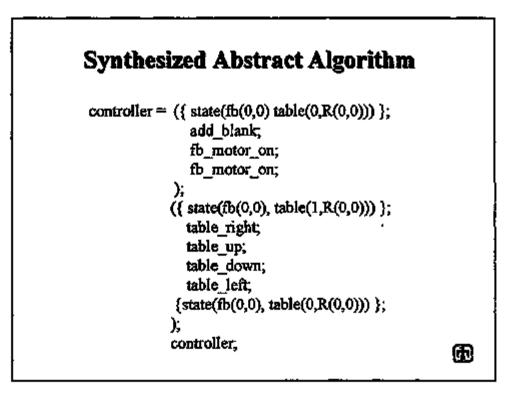
Ð

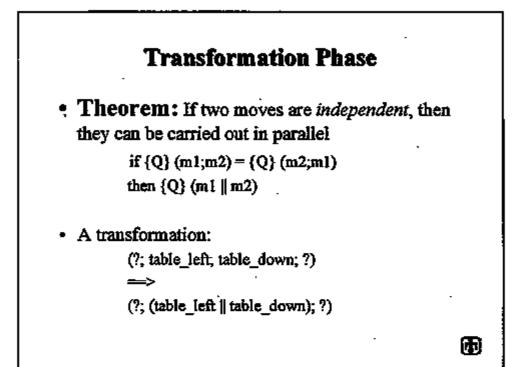
Ð

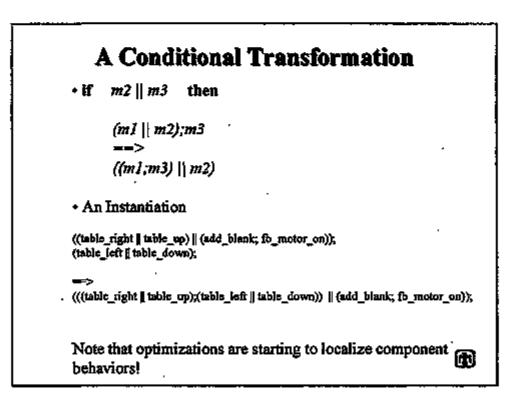
5

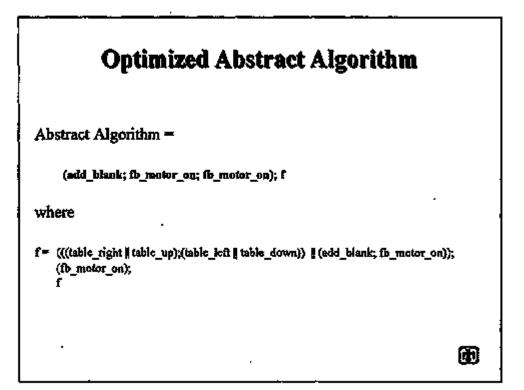
-)

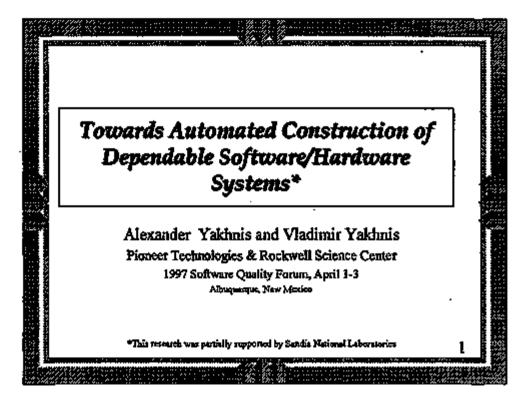


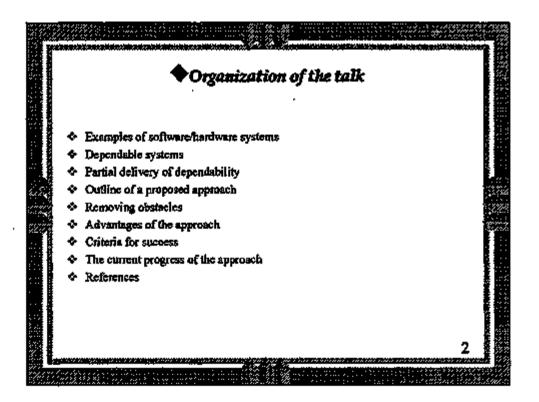


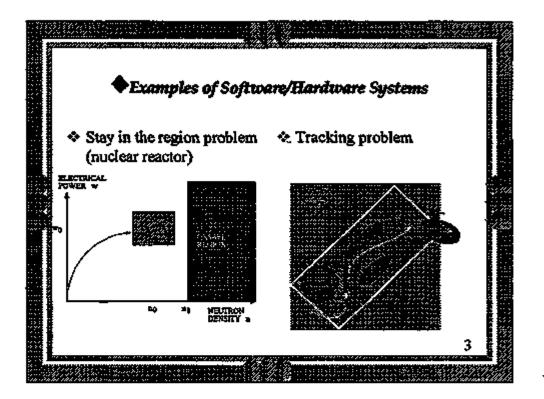


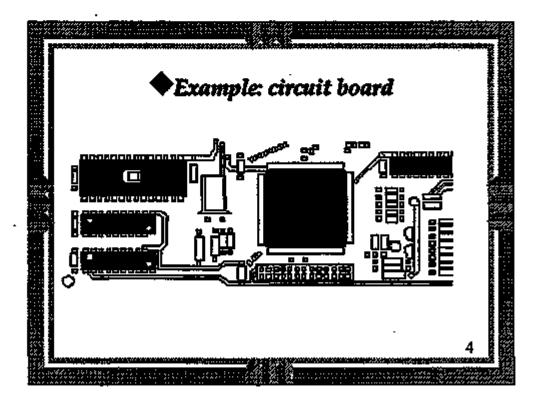






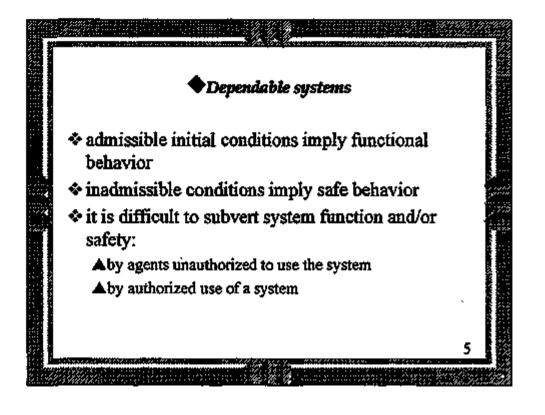


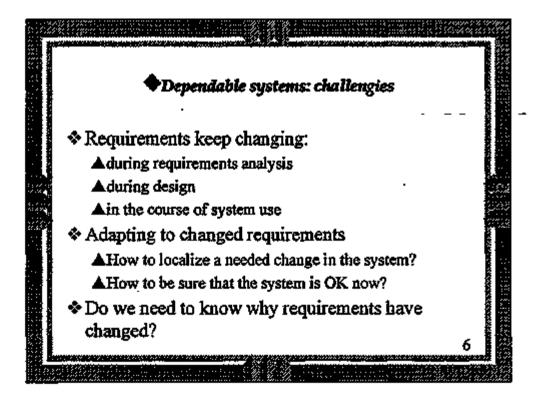


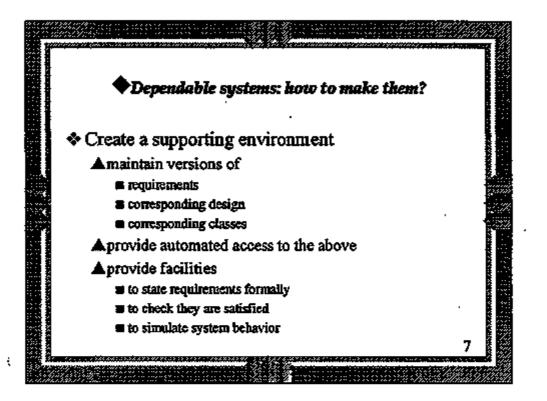


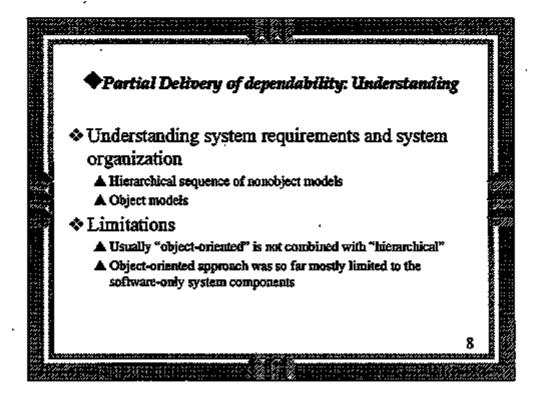
r

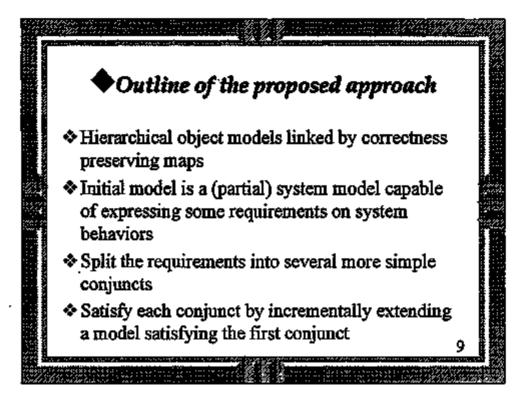
i

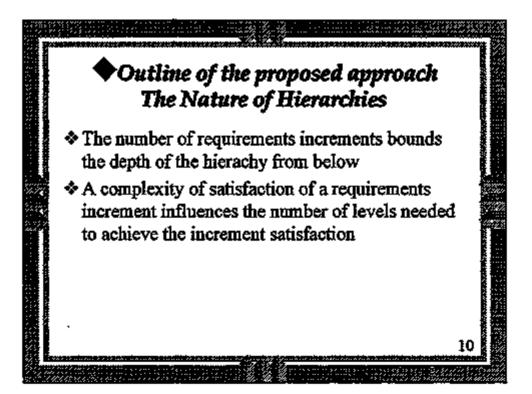


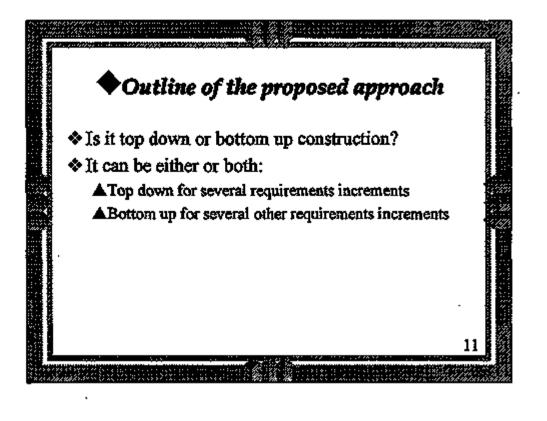


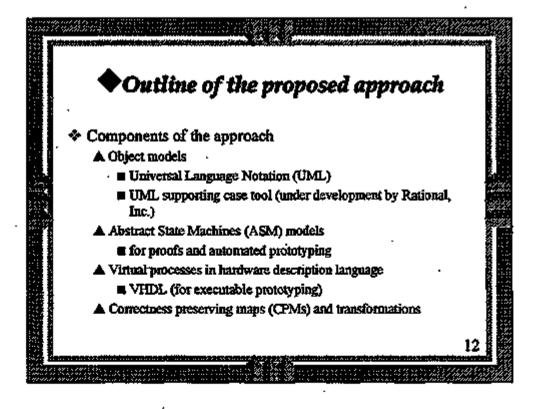




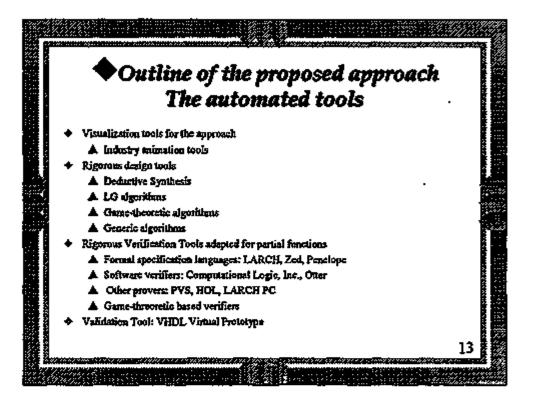


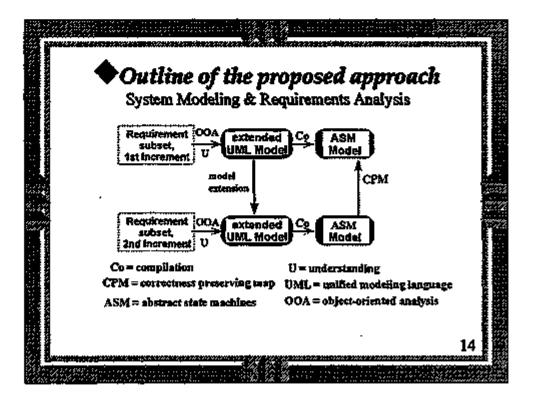


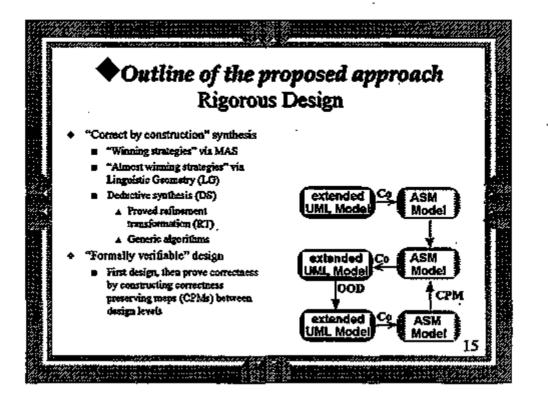


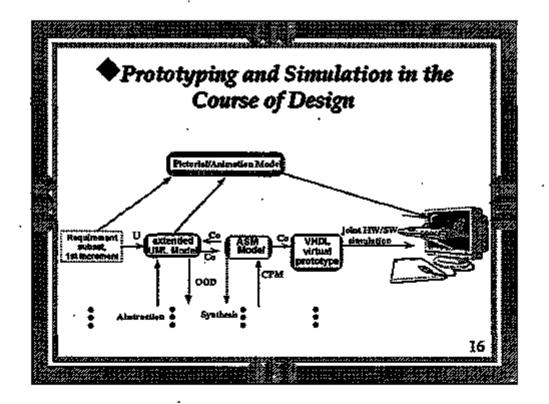


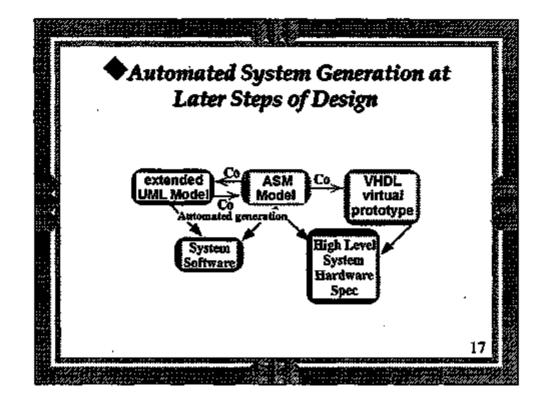
б

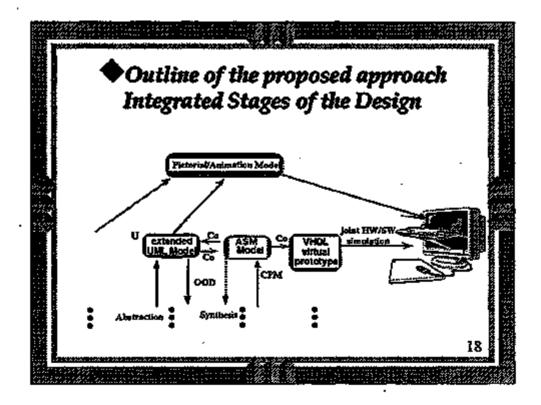


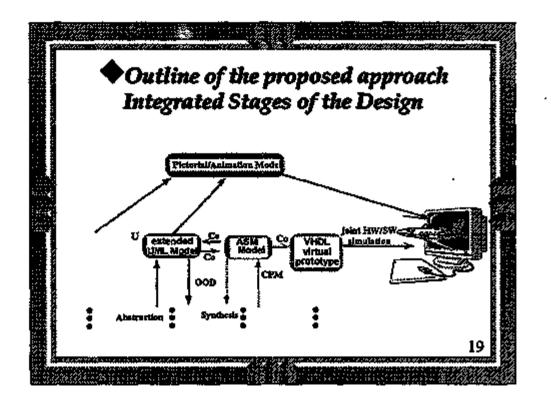


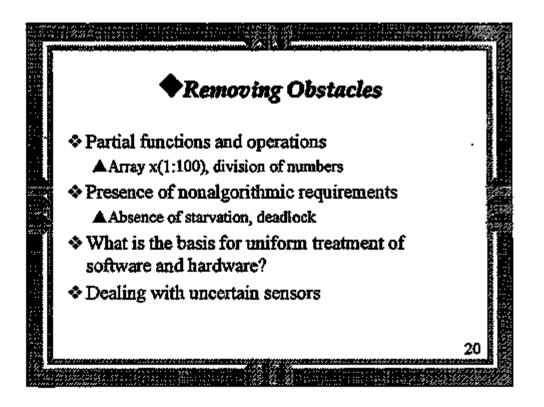


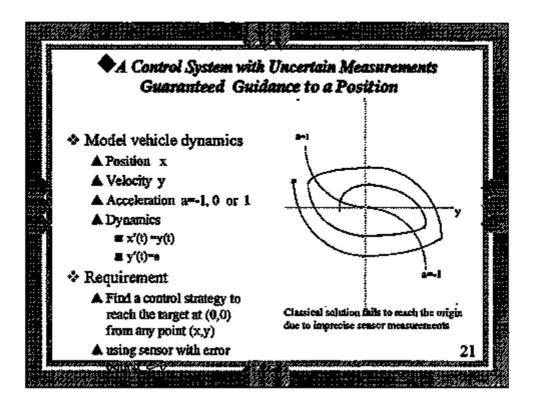


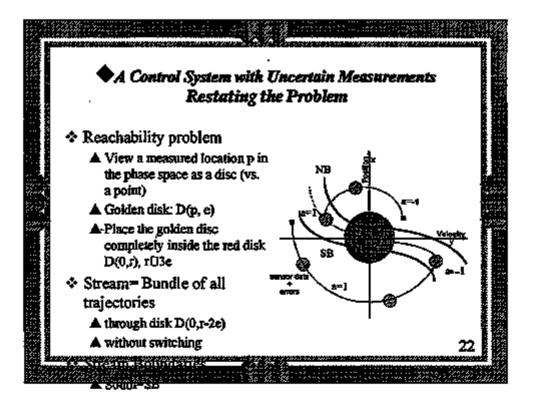


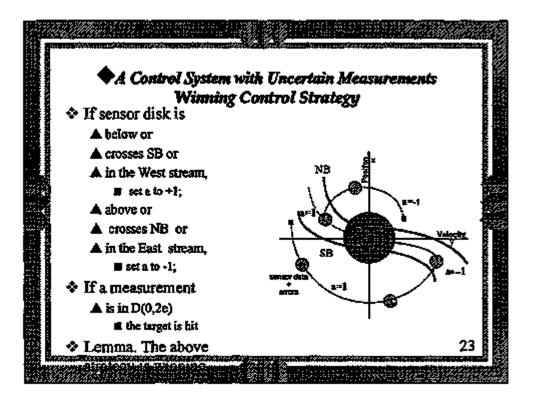


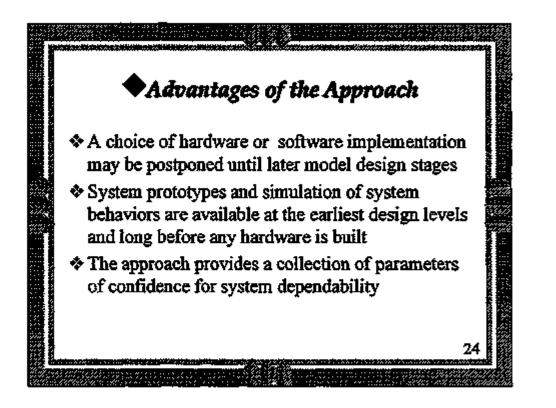


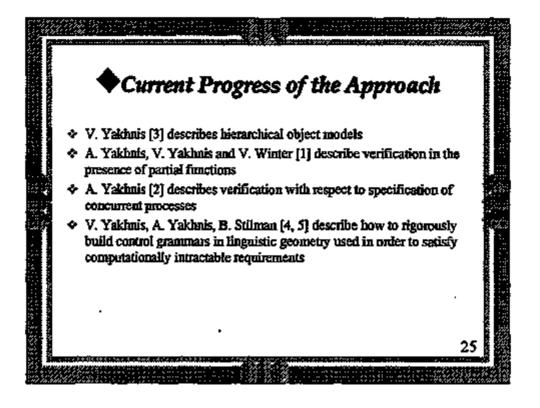


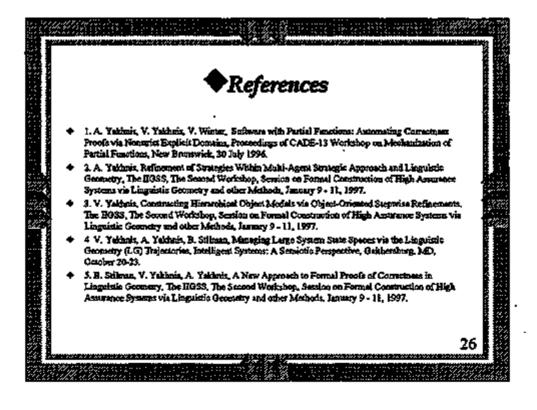


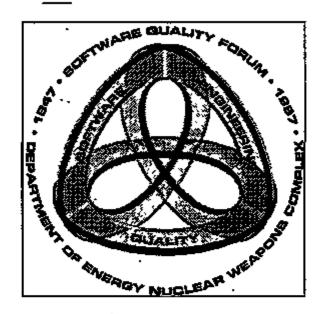








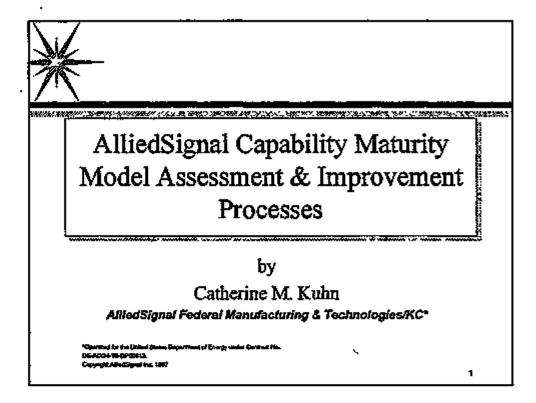


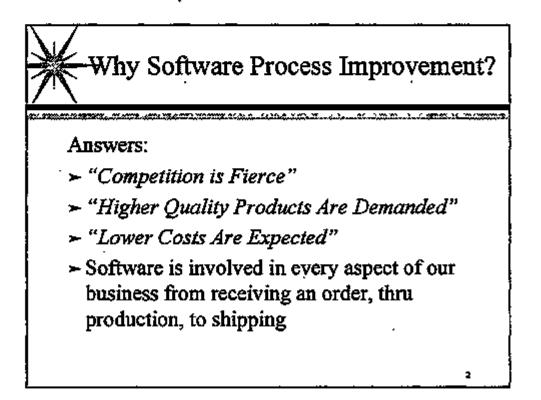


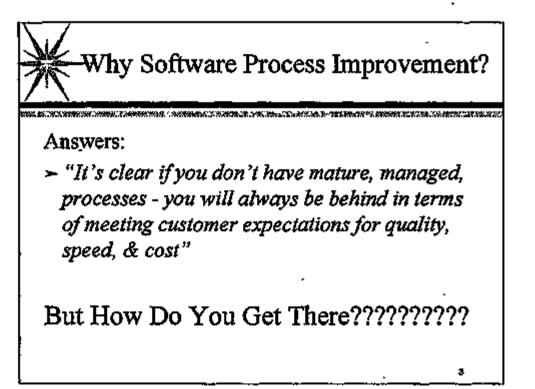
Session A4: Software Process Improvement Π

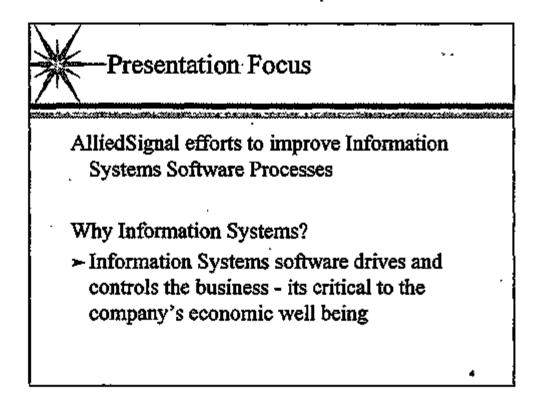
Chair John Hare AWE United Kingdom

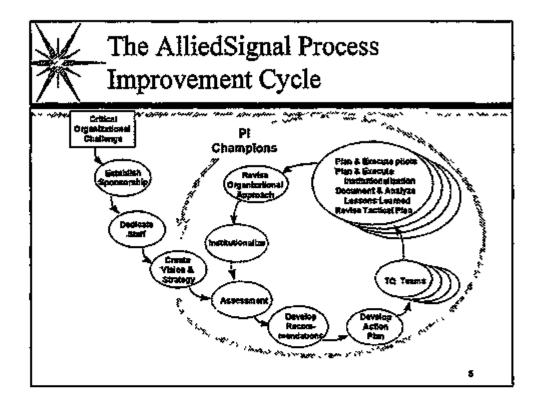
Session : Paper #	Author(s)	Title
A4:1	Cathy Kuhn AS/FM&T	AlliedSignal Capability Maturity Model Assessment & Improvement Processes
A4:2	Ann Stewart Lockheed Martin Energy Systems	Lessons Learned on Utilizing the SEI/CMM in the Federal Government Work for Others Environment
A4:3	Gail Benefield Lockheed Martin Energy Systems	"SWiM" Your Way to Software Quality

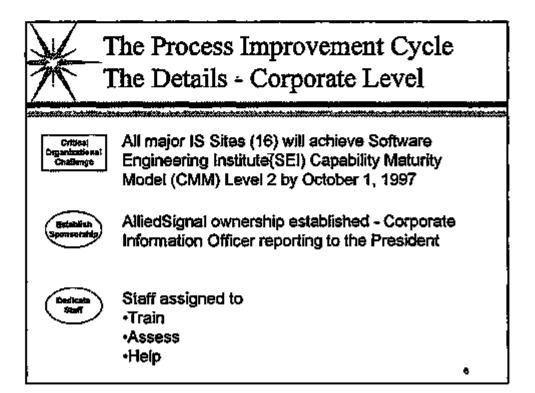


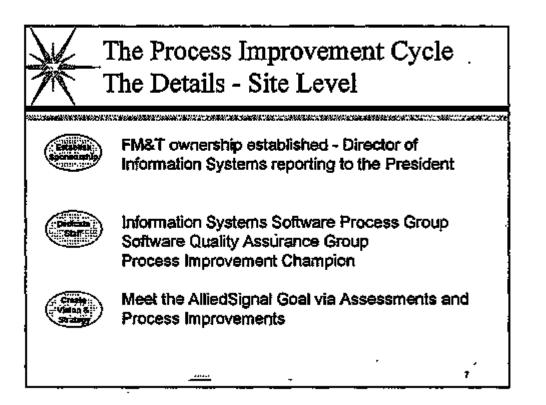


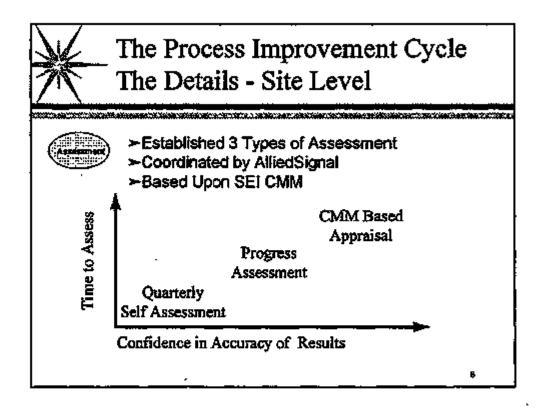


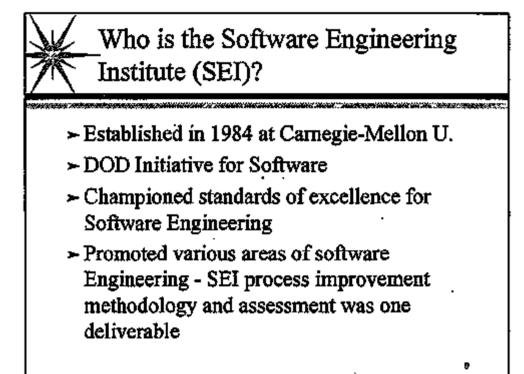


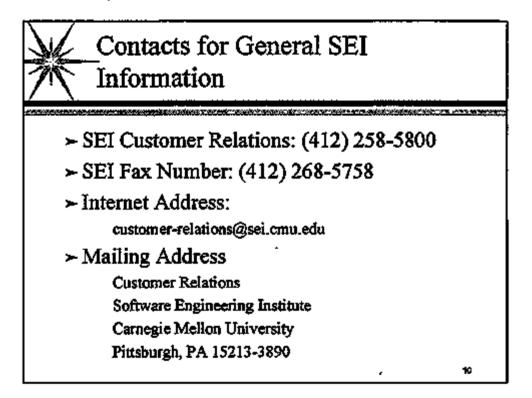




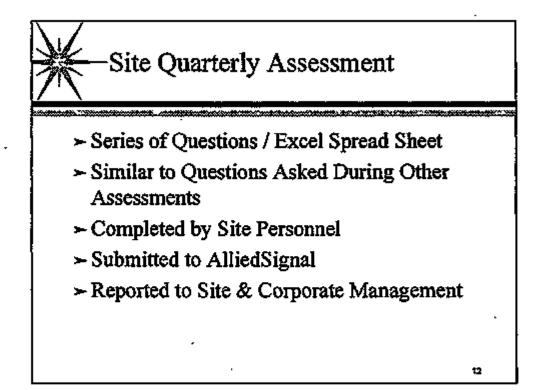


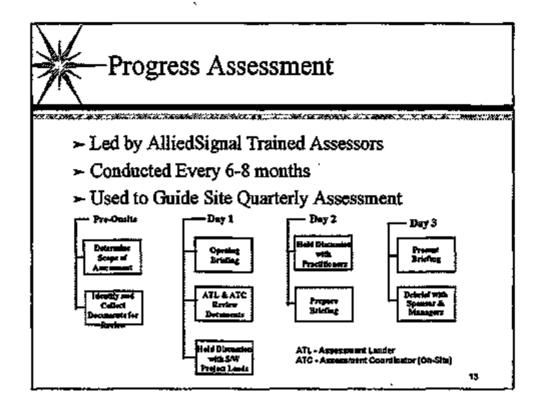


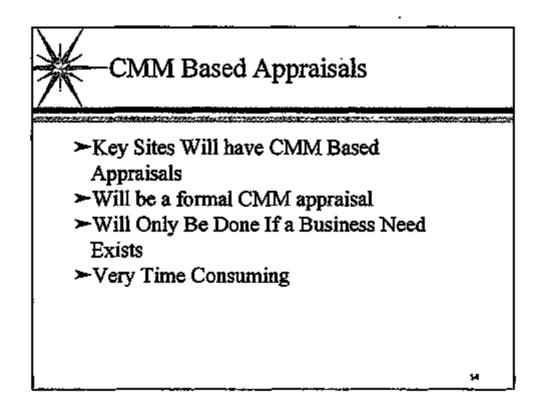


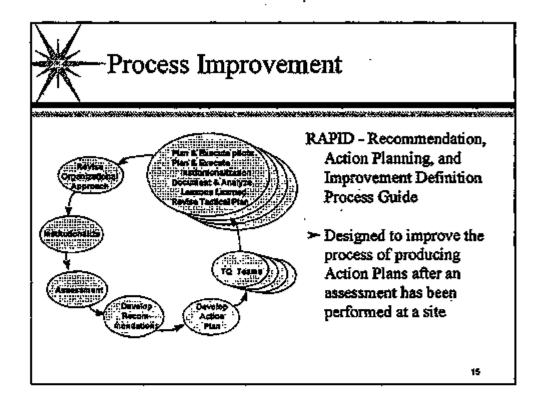


SEI's CMM Overview				
LEVEL	FOCUS	KEY PROCESS AREAS	RESULTS	
\$. Optimizing	Continuous Process Improvement	Defect Prevention Technology Change Management Proctif Change Management	Quality 3 Productions	
4. Managed	Product and Process Quality	Quantitative Process Management Software Quality Management		
• 3. Defined	Engineering Process	Organizational Process Focus Organizational Protein Definition Training Program Integrated Software Management Software Product Engineering Integroup Constituation Prer Reviews		
I. Repeatable	Project Management	Requirements Management Software Project Planning Software Project Tracking and Oversight Software Subcontract Management Software Configuration Management Software Quality Assurance		
1. Initial	Heroes	······································		

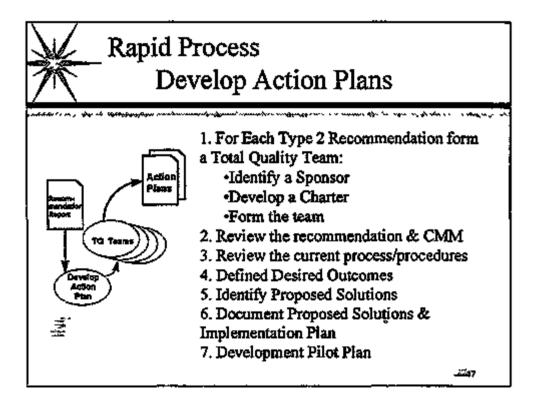


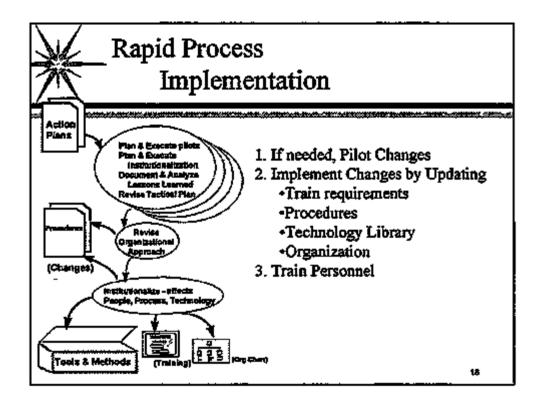


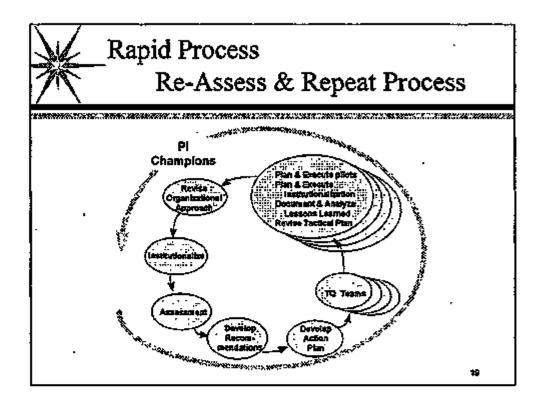


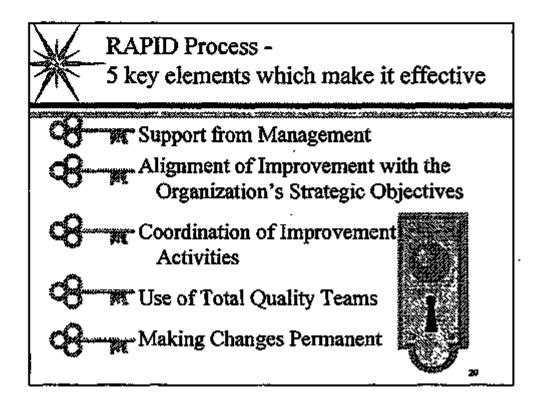


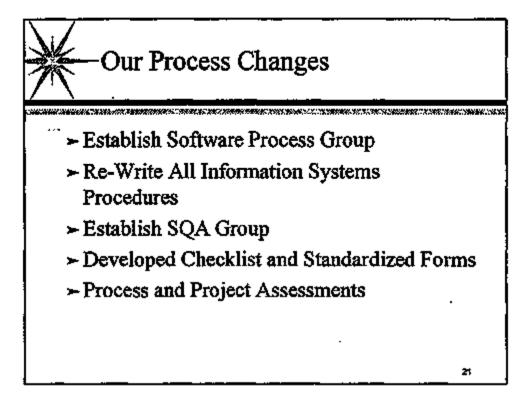
	d Process Develop Recon	nmendations		
1. Prioritize on Importance & Classify Assessment Findings				
	Type Definition	Action		
Deretop	1Obvious Solution	Assign to Individual		
	2 Needs Further	Give to		
	Investigation	Recommendation Team		
	3_Out of Scope	Give to Management		
2. For Type 2 findings Develop, Prioritize, Classify Recommendations & Report Result Type 2 Recommendations Continue on in the process				

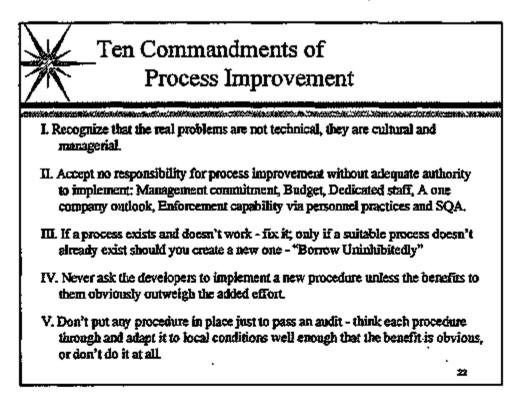


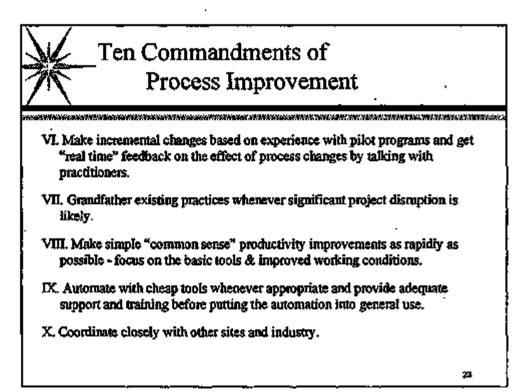


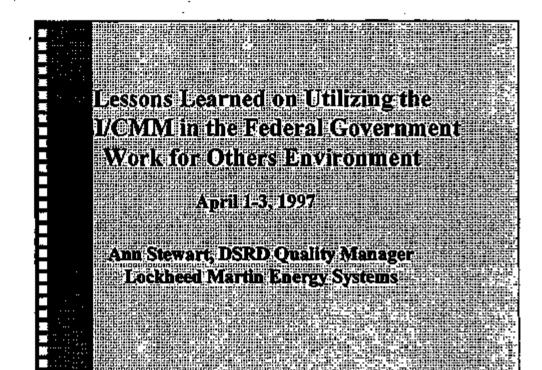








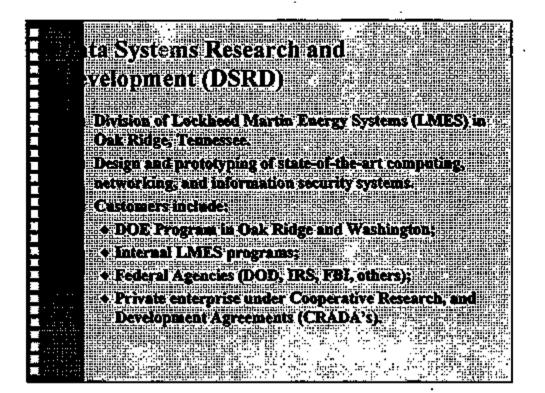


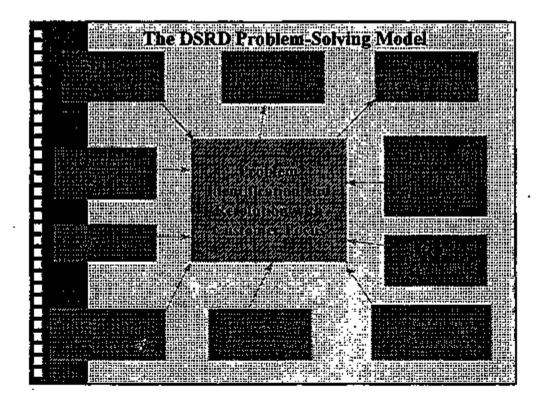


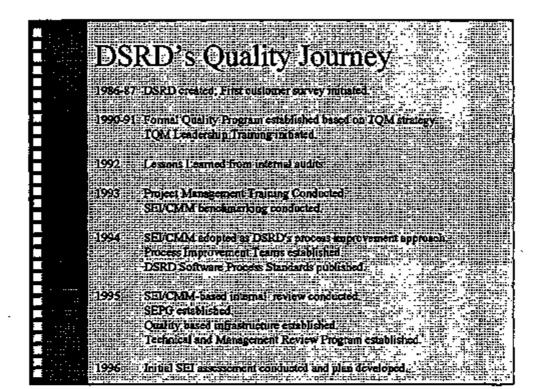
Lessons Learned on Utilizing the SEI/CMM in the Rederal Government Worksfor Others Environment

Data Systems Research and Development
 SEI/CMM

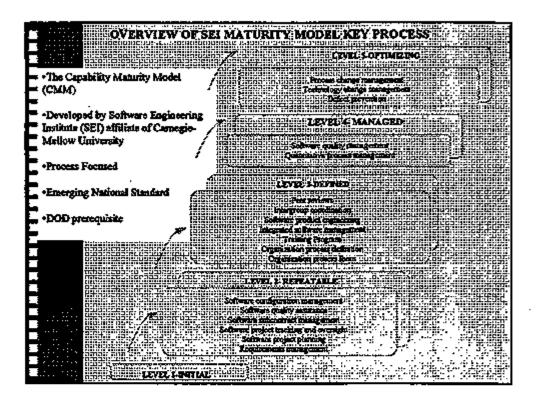
DSRD Process Improvement Approach
 Accomplishments
 Lessons Learned



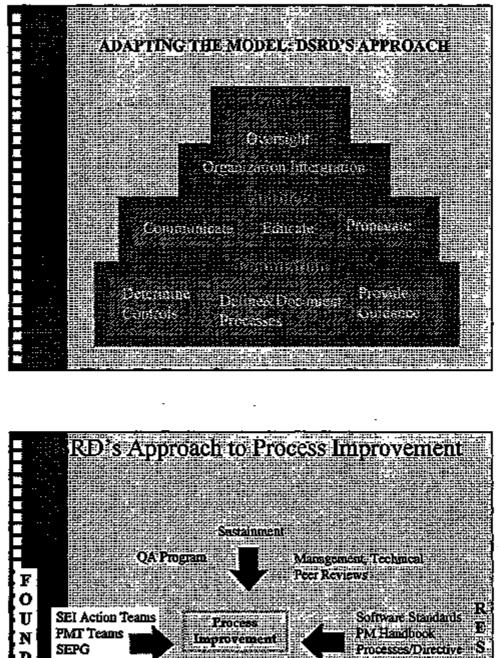




What is the SEL/CMM2 The Software Engineering Institute (SEI) is an affiliate of Carnegie-Mellon University established to address problems faced by the Software Industry. SEI developed the Capability Maturity Model (CMM) as a Framework to improve software engineering processes. CMM is an emerging national standard for evaluating capabilities of software development organizations. A 1996-97 prerequisite of Department of Defense software contracting.



Why Did DSRD Adopt This Model? A need to add structure to existing expertise to establish a more disciplined and repeatable processes. A need to retain customer focus and flexibility to commute to be responsive to customers specialized needs. Promoted by the Software Productivity Consortium (SPC). Established to promote process improvement to its member companies; Conducts SEI-based Assessments; Member companies include: Aerojet Boeing, Rockwell International, Grumman, Syscon, and Lockheed Martin.



Approach

Improvement Cycles

SEPG

D

A

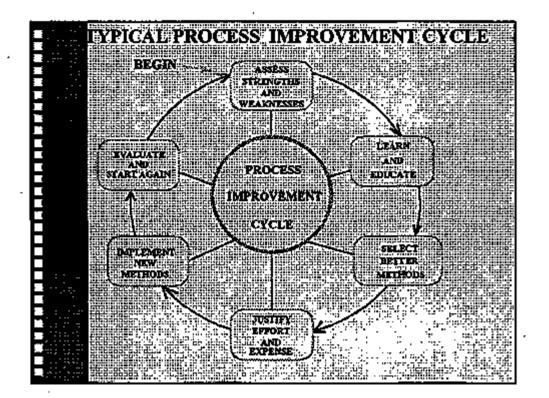
T

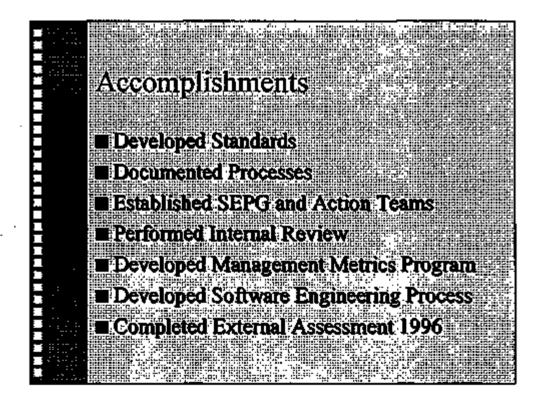
Ι

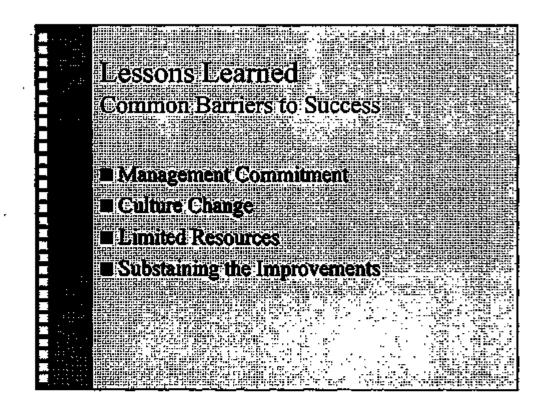
0

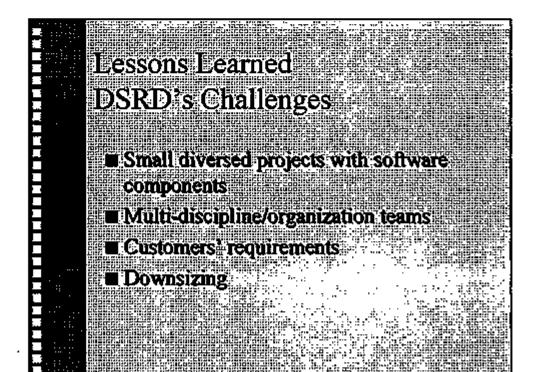
N

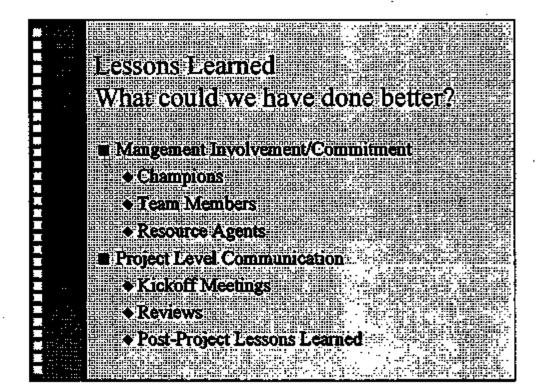
15 Action Plans U L ŝ

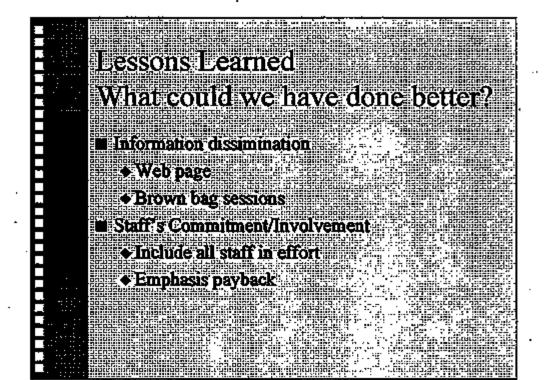




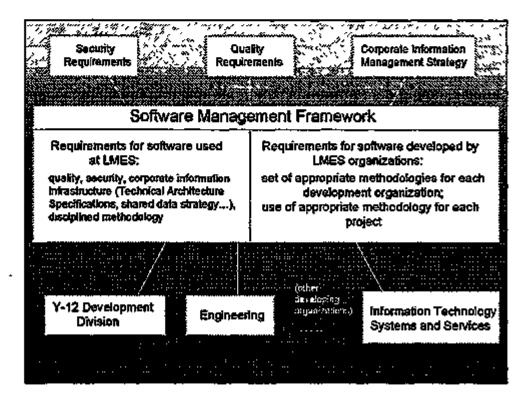


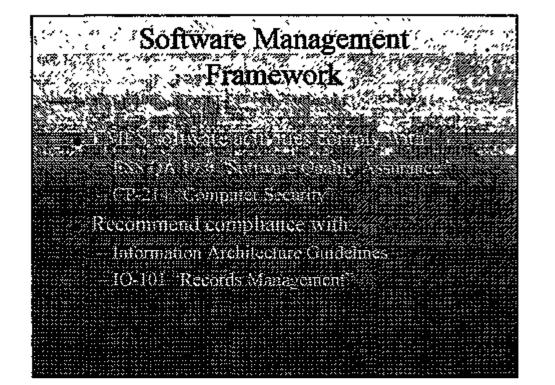


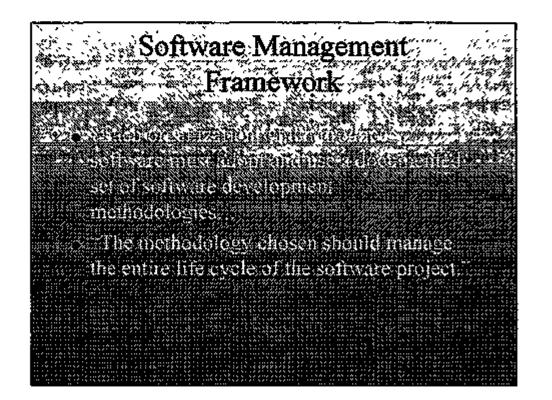




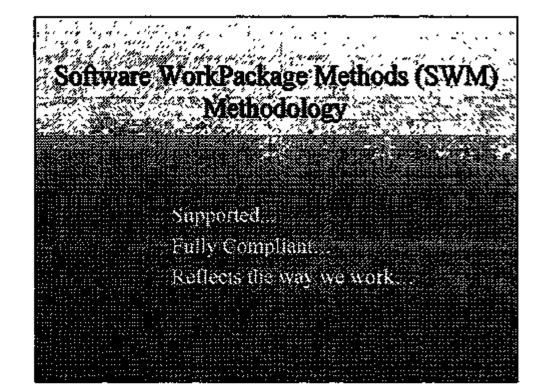


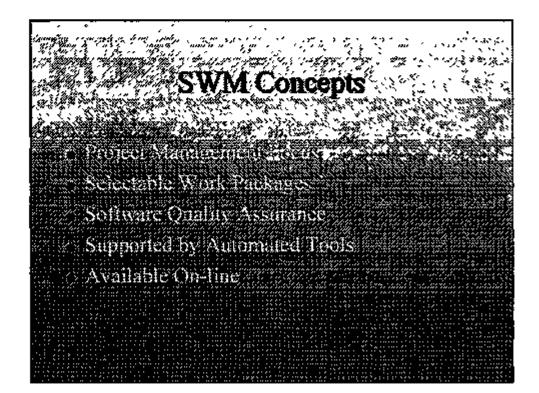


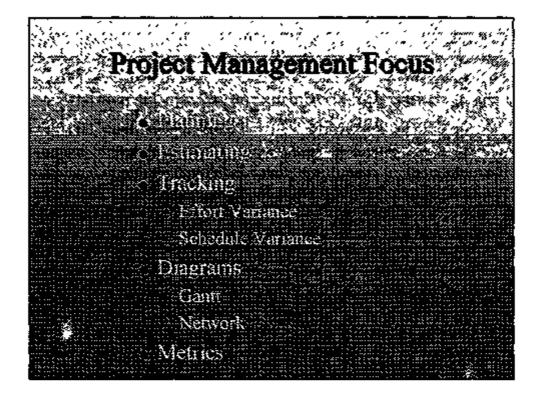


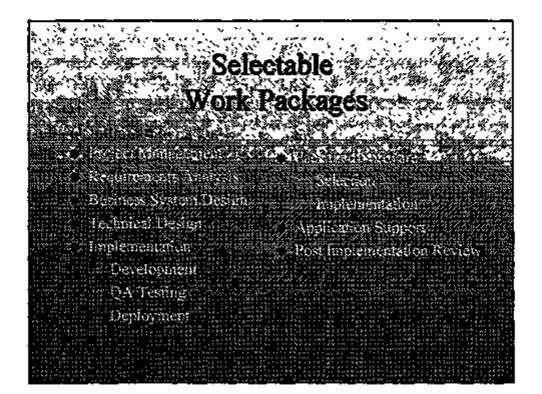


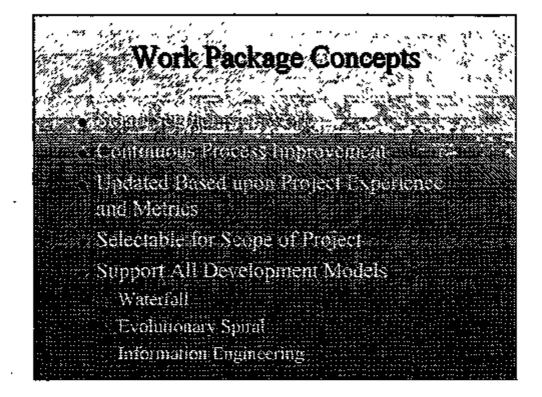
)



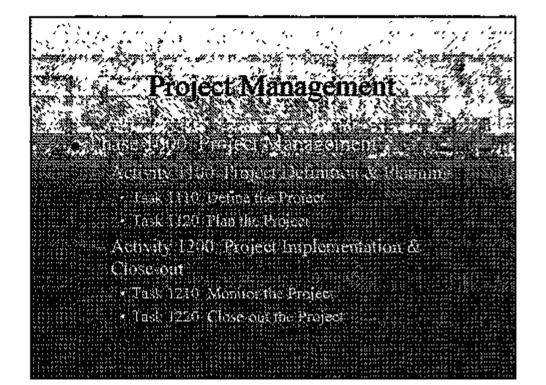


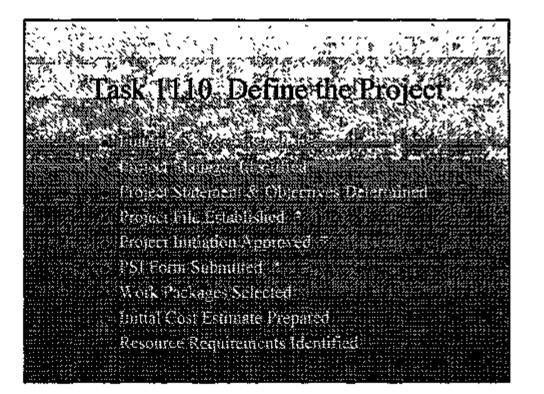




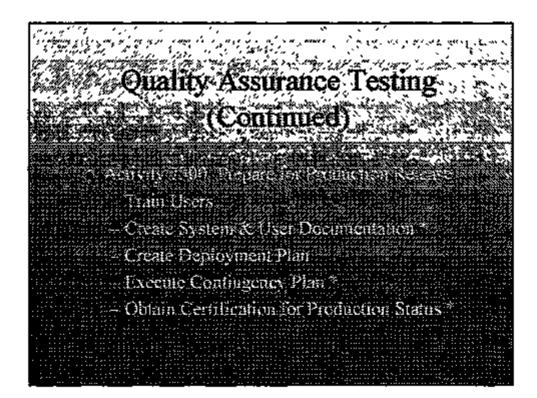


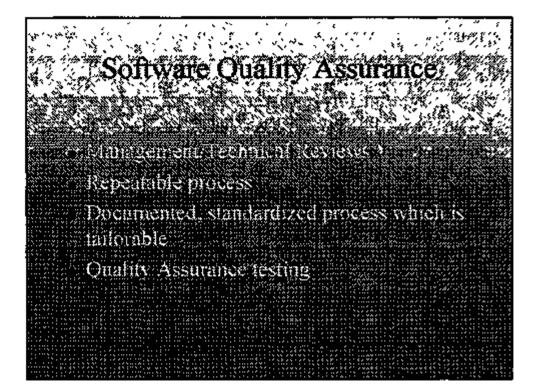


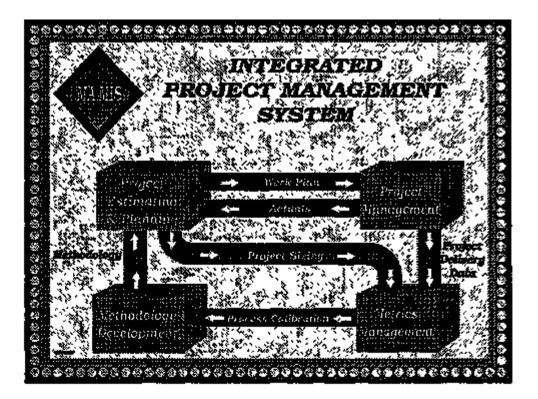


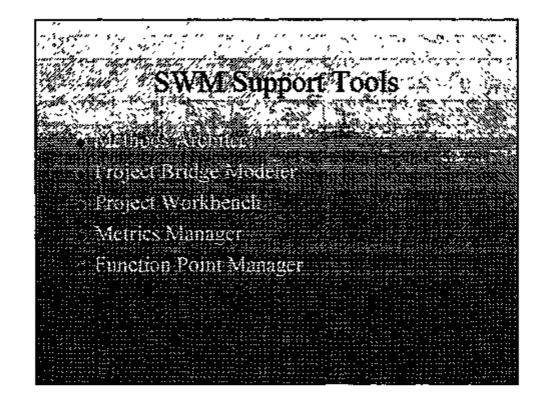


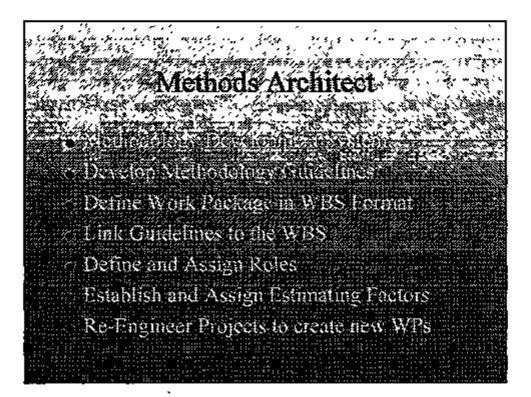
Quality Assurance Testin Create Transmion Pian Crewte User Acceptance Tost Plan Create Certification Test Plan Activity 7200 Formal Testing in QA Environment - Conduct System Test * Conduct User Acceptance fiest ** Certification Test Results Approved *

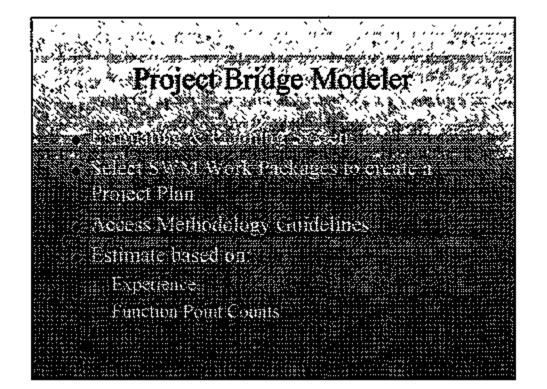


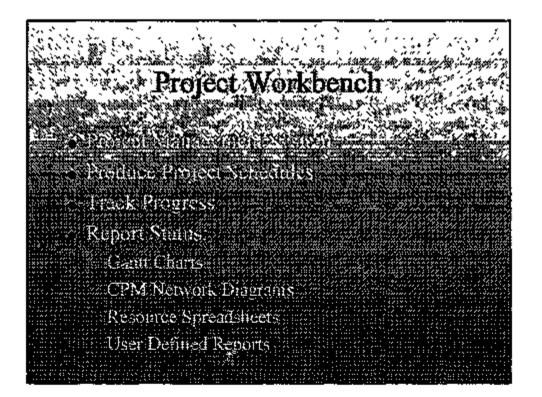


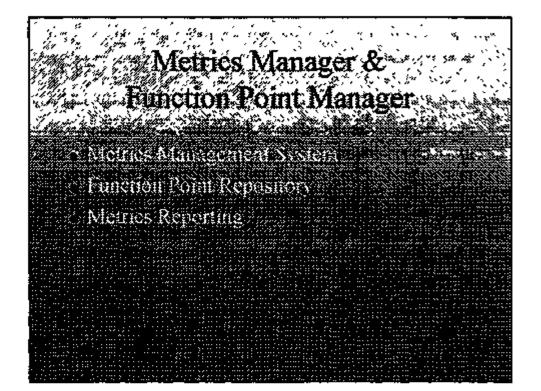


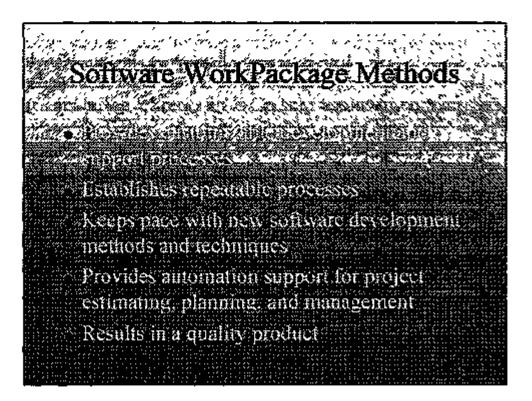


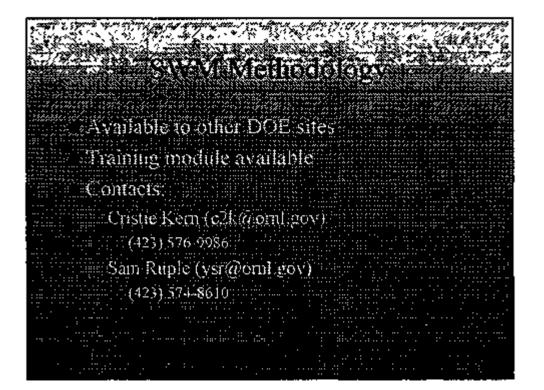




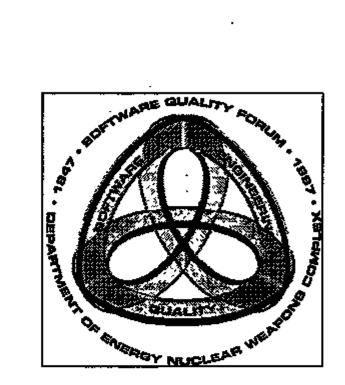








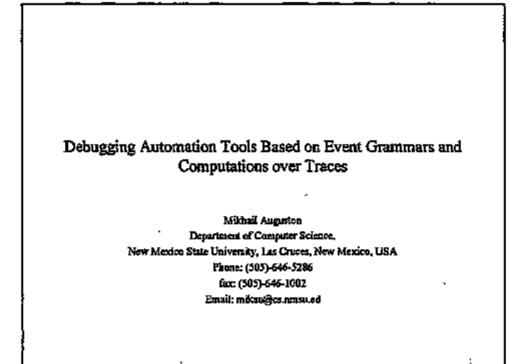
--)

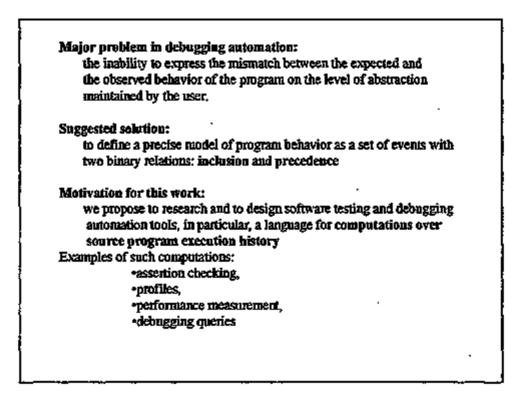


Session B4: High Integrity / Formal Methods II

Chair Larry Dalton Sandia National Laboratories

Session : Paper #	Author(s)	Title
B4:1	Mikhail Auguston New Mexico State University	Debugging Automation Tools Based on Event Grammars and Computations over Traces.
B4:2	Marie-Elena Kidd Sandía National Laboratories	A Method for Critical Software Event Execution Reliability in High Integrity Software
B4:3	John Sharp Sandia National Laboratories	Business Rule Enforcement Via Natural Language Modeling





Essential features of this approach:

The notion of an event grammar provides a precise and formal model of parallel program behavior defined as a set of partially ordered nested events

· Event attributes provide complete access to each target program's execution state

 The inclusion relation yields a hierarchy of events, assortions can be defined at appropriate level of granularity

Events can be detected by automatic source program instrumentation.

Patterns and aggregate operations on events describe computations over event traces.

Our approach is nondestructive: assertion text is separated from the source program's text

Ability to formalize universal assertions and to define debugging rules and strategies.

Events

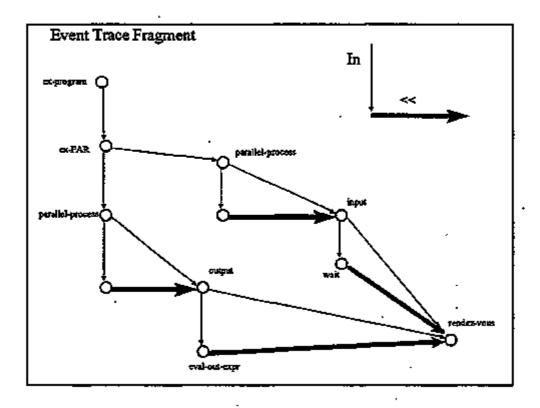
- A particular action may be performed many times, but every execution of the action is denoted by a unique event.
- · Every event is associated with a time-span that has a defined beginning and end.
- A composite event is a (partially ordered) set of other events.
- An event occurs when some action is performed in the target program execution process. For instance: a message is sent, a statement is executed, or an expression is evaluated.

2

 Each event should be detectable during the target program run time by an appropriate (automatic) instrumentation

An event gram	nar for an OCCAM subset
ex-program ::	(ex-process).
ex-process ::	(SKIP STOP ex-action ex-construction ex-instance)
ex-action ::	(ex-assignment input output)
ex-assignment ::	(sval-righthand-part destination)
eval-righthand-par	t:: (eval-expr)
destination ::	(variable array-elt)
input ::	(channel [wait] rendez-vous destination)
output ::	(channel cyal-out-copr [wait] rendsz-vous)
eval-out-expr ::	(eval-expr)
Note: input an	d output of the same message share the same rendez-yous event
ex-construction to	(ex-SEQ) ex-conditional ex-loop ex-PAR ex-ALT)
co-SEQ ::	([ex-replicator]ex-process*)
ex-conditional ::	([ex-replicator]eval-condition + ex-cond-branch)
eval-condition ::	(evai-expr)
ex-cond-branch :;	(ex-process)
ex-loop ::	(ex-one-iteration +)
	(eval-condition [ex-loop-body])
ex-loop-body ::	

```
ex-PAR ::
                   ([ex-replicator]{ parallel-process *})
parallel-process :: (ex-process)
                   ([cc-replicator] channel* (alt-wait | eval-condition)*
ex-ALT ::
                    [ex-guard] co-alternative )
ex-guard ::
                   (input)
ex-alternative ::
                   (ex-process)
                   ( variable base-expt count-expt )
ex-replicator ::
                   (eval-expr)
base-expr ::
count-expr ::
                   (eval-expr)
                  (instance-name eval-act-parameter * ex-instance-body)
ex-instance ::
eval-act-parameter :: (eval-expr destination)
ex-instance-body :: (ex-process)
                   (eval-simple-expr | eval-dyadic-expr )
eval-expr ::
eval-simple-expr :: ( constant | variable | array-elf | eval-monadic-expr )
eval-dyadic-expr :: ( eval-1st-arg eval-2nd-arg perform-bin-op )
                  (cval-cxpr)
eval-1st-arg ::
eval-2nd-arg ::
                   (eval-expr)
stray-eit ::
                   ( array-name eval-index )
eval-index ::
                   (eval-expr)
eval-monadic-expr ::
                            (eval-arg perform-mon-op)
eval-arg ::
                   (cval-expr)
```

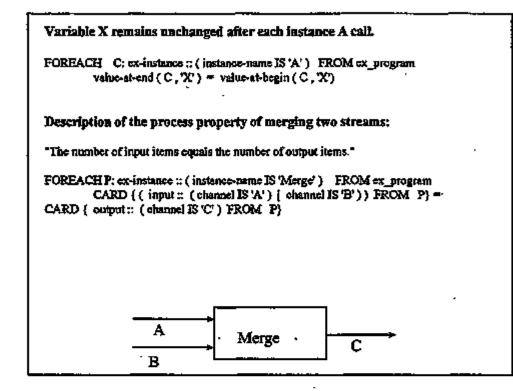


This model makes it possible to formalize assertions of the type:
"all variables in the program must be initialized before using in some expression",
"file must be opened, then the read statement is performed zero or more times and after that the close statement is executed",
"at least one variable changes its value during one loop iteration",
"after the execution of a subprogram P the value of variable X remains unchanged",

• "there is an attempt to assign values to the same variable in two parallel processes" (data race condition).

Assertion examples PAR Channel1 | Message1 Channel1 ! Message2 Dynamic constraint EXISTS Snapshot :: { O1: output, O2: output } (channel-tag(Nearest-included-channel(O1)) = -channel-tag(Nearest-included-channel(O2))) SAY 'Attempt to use channel' source-text(Nesrest-included-channel(O1)) 'in two parallel processes." source-text(Least-embracing-parallel-process(O1)) 'and' source-text(Least-embracing-parallel-process(O2)) 'in output statements' source-text(O1) 'and ' source-text(O2) 'respectively' This is an example of an universal assertion

<pre>PAR X := expr1 X := expr2 EXISTS Snapshot ::{D1: destination, D2: destination } (location (D1) = location (D2)) SAY ('Attempt to assign to the same memory location' source-text(D1) 'and' source-text(D2) 'in two parallel processes:' source-text(Least-embracing-parallel-process(D1)) 'and' source-text(Least-embracing-parallel-process(D2)))</pre>	X :=	expr1	
X := expr2 EXISTS Snapshot ::{D1: destination, D2: destination } (location (D1) = location (D2)) SAY ('Attempt to assign to the same memory location' source-text(D1) 'and' source-text(D2) ('in two parallel processes:' source-text(Least-ombracing-parallel-process(D1)) 'and'		expr1	•
X := expr2 EXISTS Snapshot ::{D1: destination, D2: destination } (location (D1) = location (D2)) SAY ('Attempt to assign to the same memory location' source-text(D1) 'and ' source-text(D2) ('in two parallel processes' source-text(Least-ombracing-parallel-process(D1)) 'and'			
<pre>(location (D1) = location (D2)) SAY ('Attempt to assign to the same memory location' source-text(D1) 'and ' source-text(D2)</pre>	X :=	expr2	
source-text(D1) 'and ' source-text(D2) < 'in two parallel processes.' source-text(Least-ombracing-parallel-process(D1)) 'and'		Snapshot ::{D1: destination, D2: destination } (location (D1) = location (D2) }	• .
source-text(Least-embracing-parallel-process(D2)))		source-text(D1) 'and ' source-text(D2) 'in two parallel processes." source-text(Least-embracing-parallel-process(D1))	
- Yet another example of an universal debugging rule			



SAY	'Total time is'
	+/ { ABC: ex-instance ::(instance-name IS 'ABC') FROM ex-program APPLY duration(ABC) }
Samp	es of possible profile request
SAY	"Total number of parallel processes executed is" CARD { ALL parallel-process FROM ex-program}

References.

[Auguston, Fritzson 93] M.Auguston, P.Fritzson, PARFORMAN - an Assertion Language for Specifying Behavior when Debugging Parallel Applications, in Proceedings of the Euromicro Workshop on Parallel and Distributed Processing, Gran Canaria, January 27-29, 1993, IEEE Computer Society Press.

- [Auguston 94] Auguston M., A Language for Debugging Automation, in Proceedings of 6th Intl Conference on Software Engineering and Knowledge Engineering SEKE'94, Junuals, 1994, pp. 102-115
- [Fritzson, Auguston, Shahmehri 94] P. Fritzson, M. Auguston, N. Shahmehri: Using Associations in Doclarative and Operational Models for Automated Debugging, Journal of Systems and Software, v.25, No 3, June 1994, pp. 223-239.
- [Auguston 95] Mikhail Auguston, "Program Behavior Model Based on Event Grammar and its Application for Debugging Automation", in Proceedings of the 2nd International Workshop on Automated and Algorithmic Debugging AADEBUG'95, Saint-Malo, France, May 1995.
- [Auguston, Fritzson 96] M. Auguston, P. Fritzson, PARFORMAN an Assertion Language for Specifying Behavior when Debugging Parallel Applications, International Journal on Software Engineering and Knowledge Engineering, Vol. 6, No 4, 1996, pp.609-640.

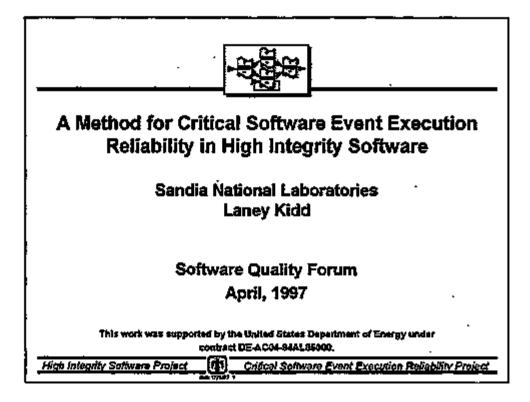
Experiments with the prototype implementation of PASCAL assertion checker have demonstrated some interesting features:

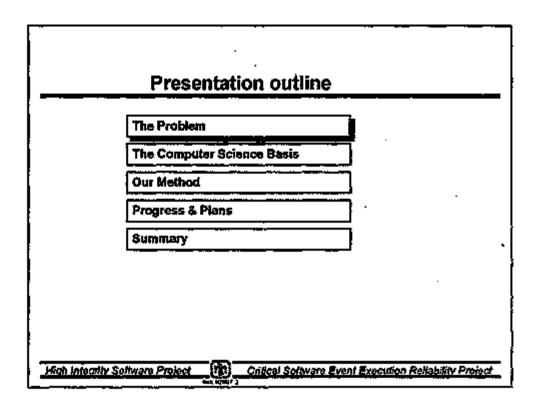
 different kinds of dynamic analysis can be described as an appropriate computations over the trace, e.g. debugging queries, assertion checking, profile measurement,

 computations over traces may provide values which otherwise can not be found in program states,

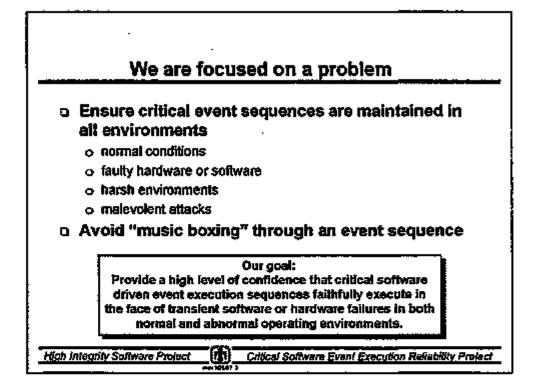
informative and readable messages can be generated,

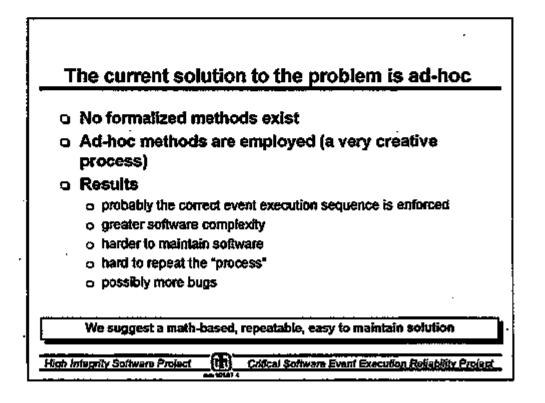
 universal assertions and debugging rules can be presented as computations over traces.

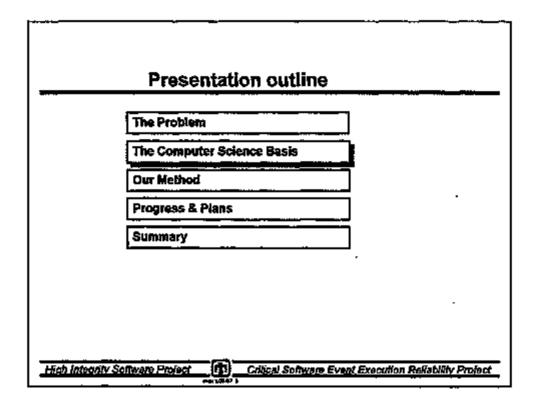


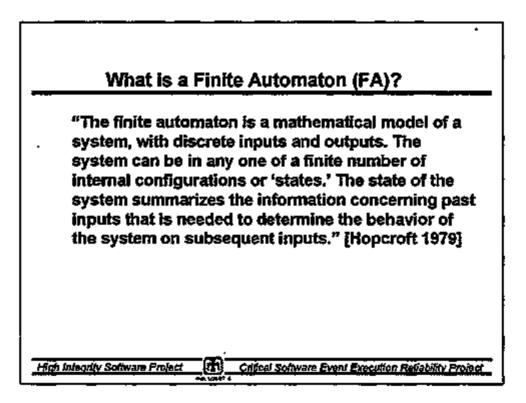


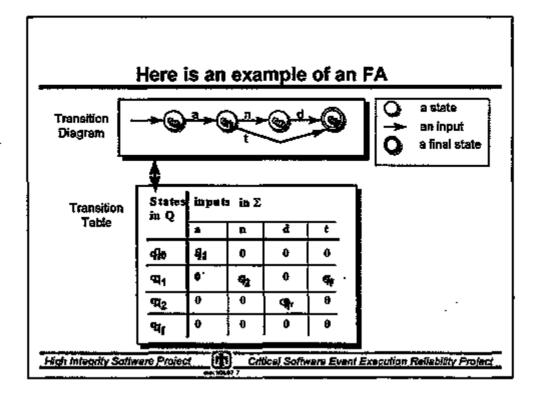
I

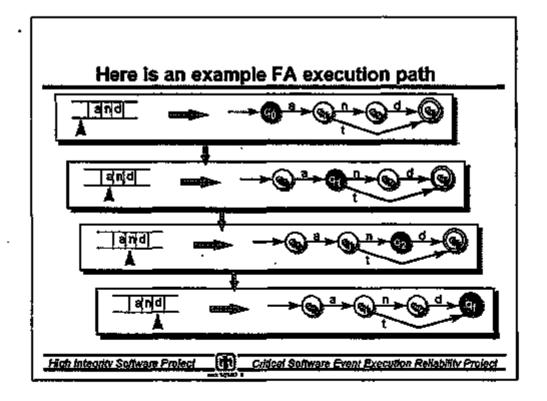








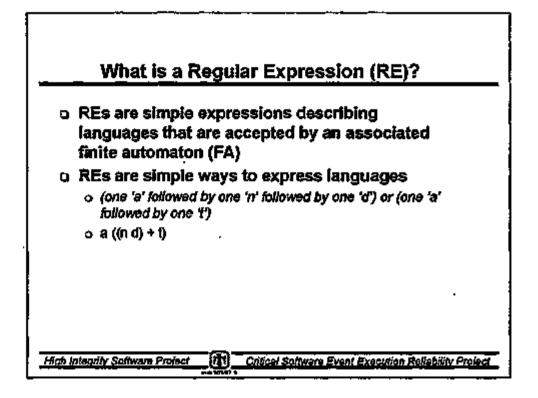




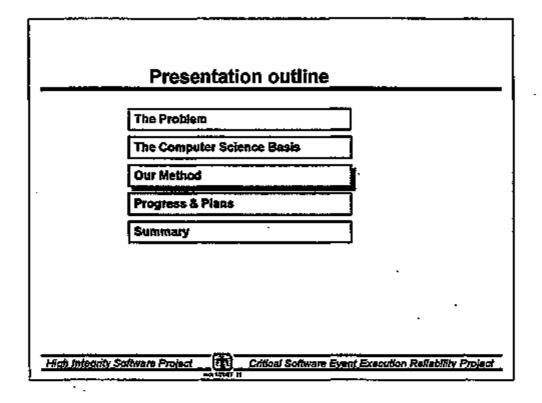
4.

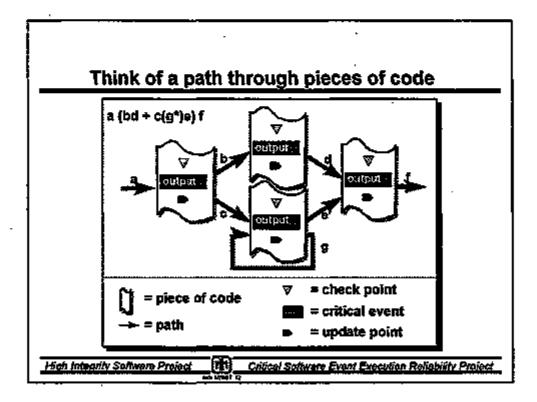
1

ı.



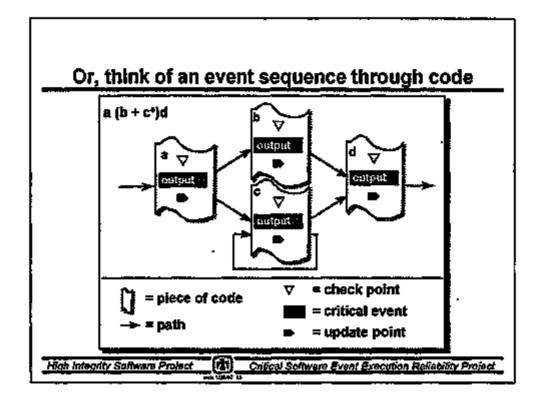
		Regular E	What is t xpression (he (RE) notation?
	0 A =	and B be so {b, c} {all, cat, at}	ets of input s	ymbols
04	Relat	ions I	Meaning	Example
	AB	Concatenation	A followed by B	birth infancy childhood adulthood A B = (ball, boat, bat, call, coat, cat) b oat = (boat)
	A+B	Selection	A or B	dog + cat + repôle + fisia A + B = {b, c, all, cat, at} b + cat = {b, cst}
	A*	Kleene Closure	0 or more	automobiles" a" = {c, a, aa, aaa,}
	A*	Positive Closure	1 or more	doctors-on-duty*
High	integni)	Software Project	Critical S	oftware Event Execution Reliability Project

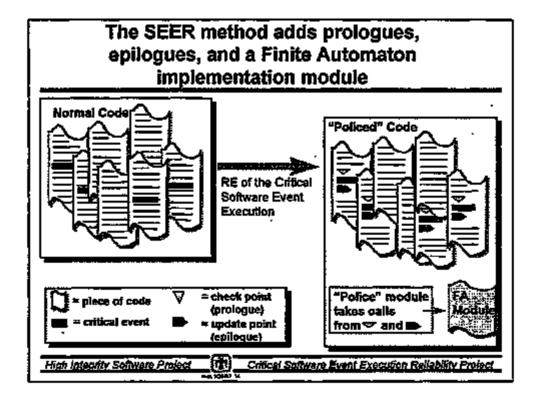


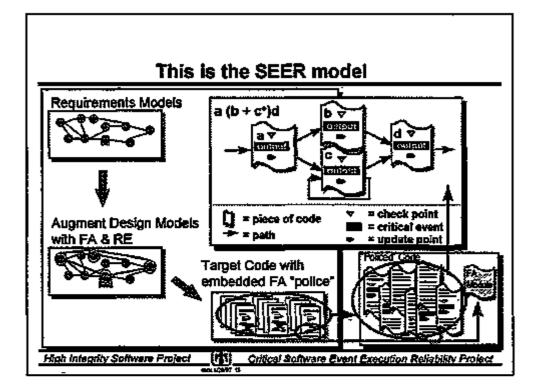


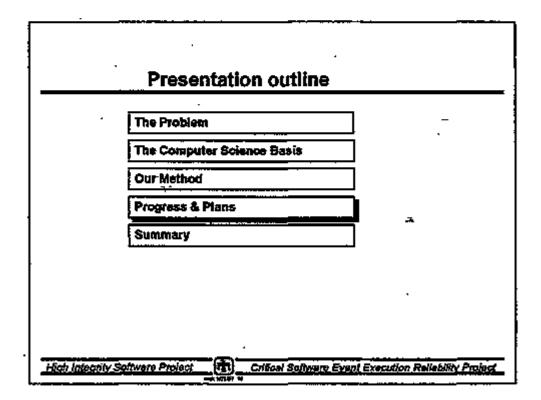
б

}









٤

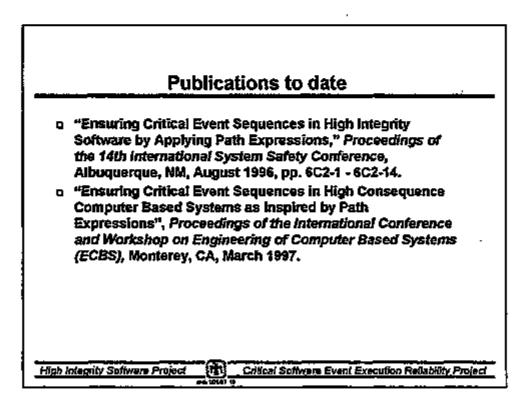
L

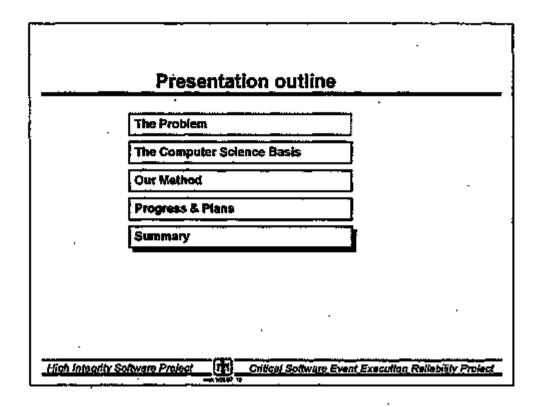
ı.

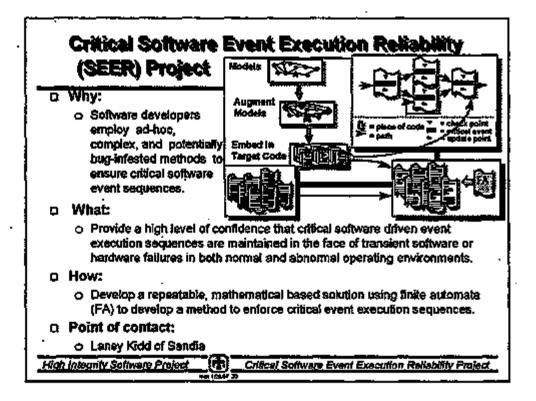
į.

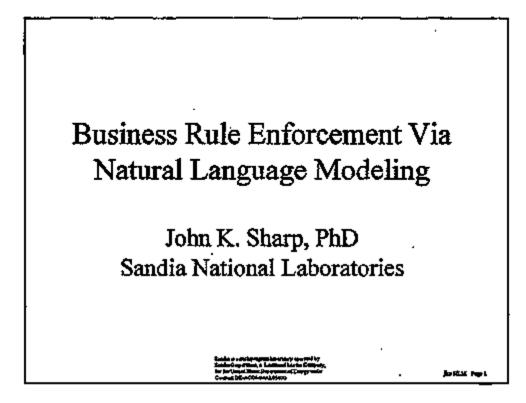
	Our pr	ogres	ss and pla	ins at a	glance	3
FY96			FY97		F¥98	
Researc		l r		1		13
& plann	ing	1		, 		1 1
	methods to ensure		ocessor <u>fault detection</u> e software critical event on Mathematics and		Apply methods to <u>fault correction</u> and distributed environments	
	Create ini method	tial	Benchmark method	Automate FA modul creation	• •	
		Create				
i i	f f	demo		i		1 1
	5 <u>7</u>	'	1 t I	· •		

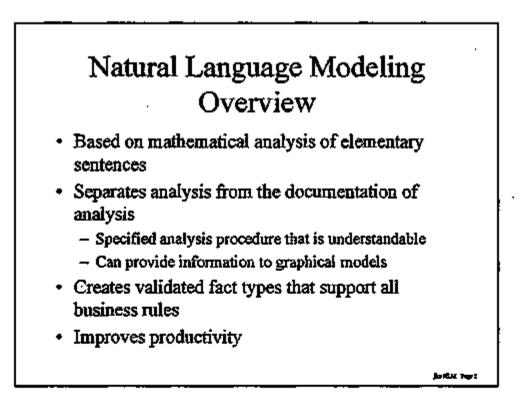
í



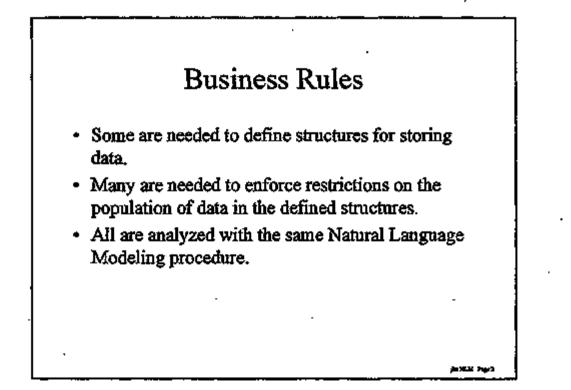


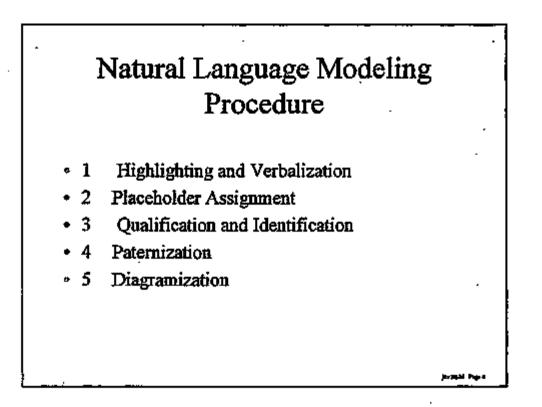


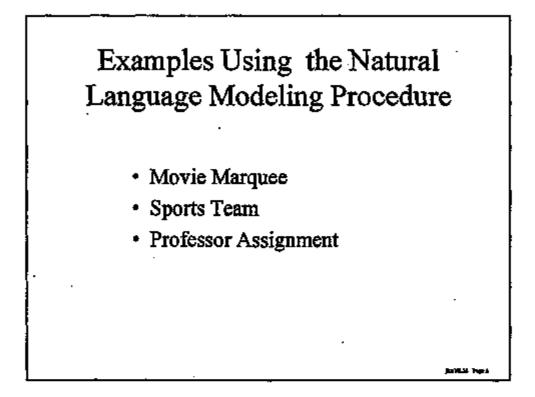


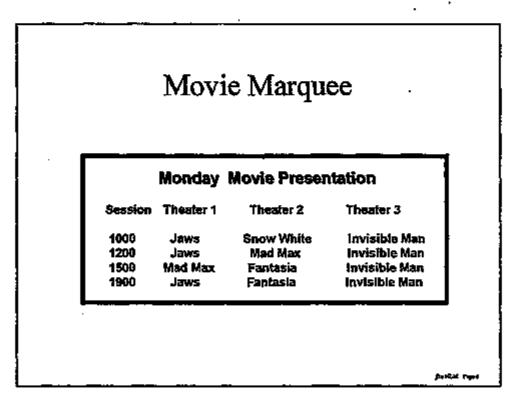


I



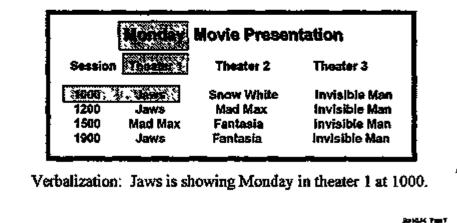


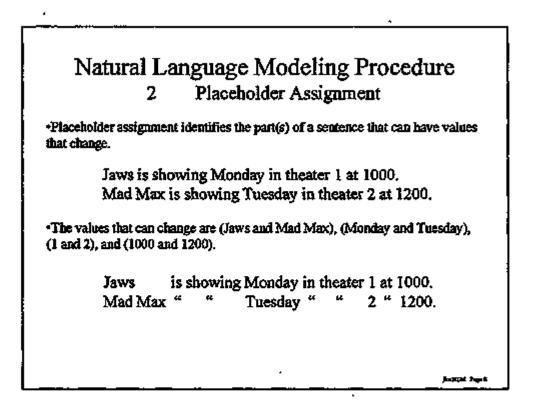




Natural Language Modeling Procedure I Highlighting and Verbalization

•Verbalization and highlighting is done by highlighting a limited example of information in the subject area and asking the subject matter expert to create a sentence.





Natural Language Modeling Procedure 3 Qualification and Identification

Jaws is showing Monday in theater 1 at 1000.

•The sentence is now tested to determine if a valid fact type can be qualified. Potential Fact Type:

<MovieName> is showing <Day> in theater <TheaterNumber> at <Time>.

Jav	/ 5	Monday	1	1000	Allowed?
8001	her	Monday	1	1000	N
Jav	¥5	another	1	1000	Y
Jav	VS	Monday	another	1000	N
Jav	vs	Monday	1	another	Y

Question: Given that fact instance "Jaws is showing Monday in theater 1 at 1000." is true, is it allowed for another valid Movie [for example "Mad Max"] to exist such that the fact instance "Mad Max is showing Monday in theater 1 at 1000." is true? Answer=No

Joséfie August

Natural Language Modeling Procedure Qualification and Identification (cont.) 3 •The sentence analysis produced two "N" answers so the corresponding objects must be analyzed together in a sentence to determine if they are independent. Potential Fact Type: <MovieName> is showing in theater <TheaterNumber>. Jaws 1 Allowed? another Ŷ 1 Y Jaws another Question: Given that fact instance "Jaws is showing in theater 1." is true, is it allowed for another valid Movie [for example "Mad Max"] to exist such that the fact instance "Mad Max is showing in theater 1." is true? Answer=Yes Result: Movie and theater are independent of each other, so two sentences must be created from the two previous "Y" answers and either movie or theater. Janik Maple

Natural Language Modeling Procedure 3 Qualification and Identification (cont.)

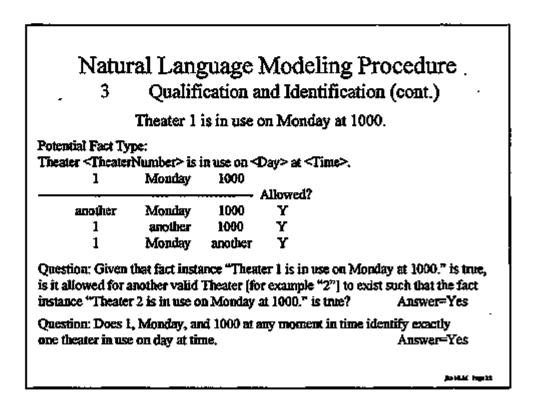
Jaws is showing Monday at 1000.

Potential Fact Type:			
MovieName> is sh	owing <day></day>	≥at <time>,</time>	,
Jaws	Monday	1000	
	•		Allowed?
another	Monday	1000	Y
Jaws	another	1000	Y
Jaws	Monday	another	Y

Question: Given that fact instance "Jaws is showing Monday at 1000." is true, is it allowed for another valid Movie [for example "Mad Max"] to exist such that the fact instance "Mad Max is showing Monday at 1000." is true? Answer=Yes

Question: Does Jaws, Monday, and 1000 at any moment in time identify exactly one movie showing on day at time. Answer=Yes

الاجهال الأراكان



Natural Language Modeling Procedure 4 Paternization

•Paternization is the specification of the general fact type that can be populated with instances.

FTI: <MovieName> is showing <Day> in theater <TheaterNumber> at <Time>.

Natural Language Modeling Procedure 5 Diagramization

-Diagramization presents a relational diagram that can be populated with instances and read using the associated fact type(s).

Movie_Day_Time

Movie Name	Day	Theater Number	Time
Jaws	Monday	1	1000
Snow White	Monday	2	1000
Mad Max	Tuesday	1	1200

Januar Papela

JANUA PARTS

Movie Marquee Enforced Business Rules

- 1 Only one movie can be shown at a time in a theater.
- 2 Only one copy of a video tape will be leased at any time.

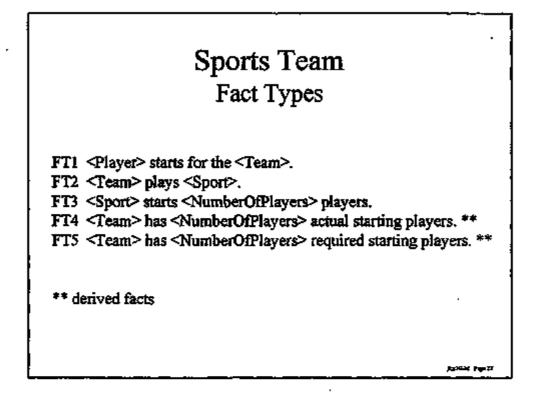
Sports Team Problem Statement

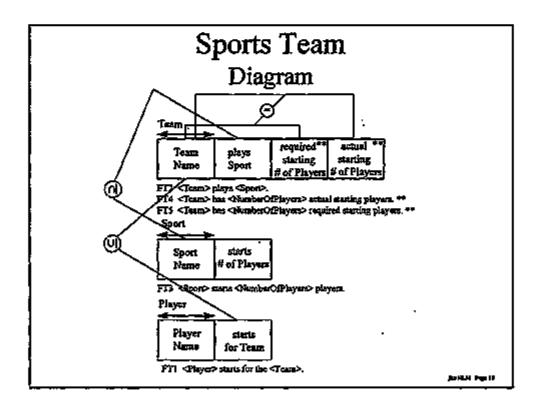
A player can start for only one team. A team plays only one sport. A sport has a required number of starting players. A team must start the number of players required for the sport the team plays.

Janual Ine Id

In the later of th

}





Sports Team Enforced Business Rules

A player can start for only one team.

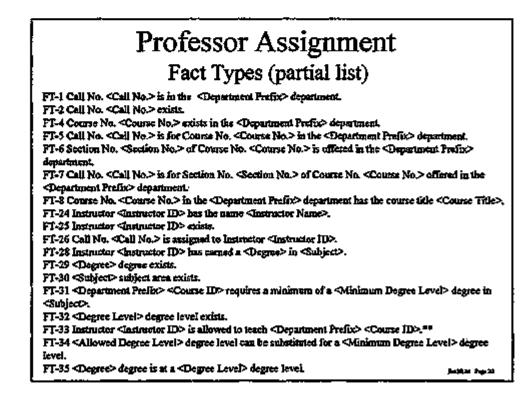
A team plays only one sport.

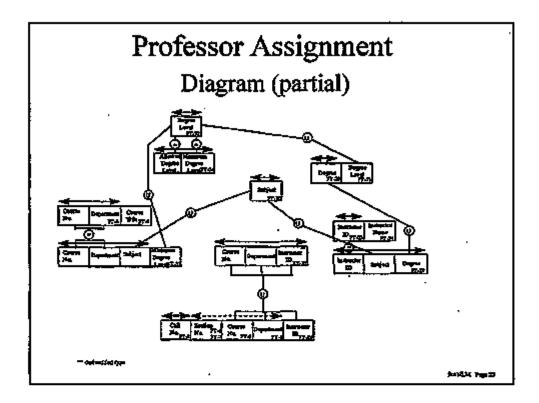
A sport has a required number of starting players. A team must start the number of players required for the sport the team plays.

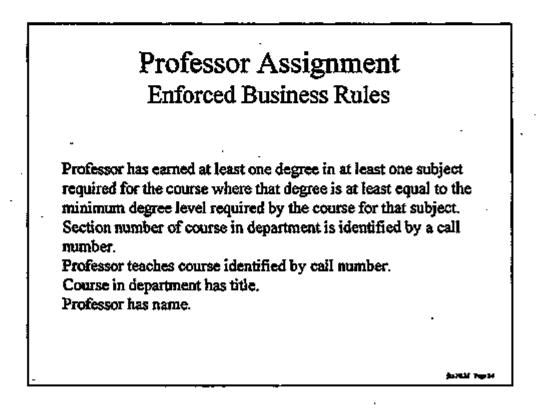
Justia Ny H

Professor Assignment General Requirements * Course ID exists in the database **(1)** (2) Professor ID exists in the database, (3) Professor has earned at least one degree in at least one subject required for the course where that degree is at least equal to the minimum degree level required by the course for that subject. (4) Section ID exists in the database. Section is for the designated course. (5) Section is not already assigned to be taught by another Professor. 6) (7) Professor is not already teaching four sections. (8) Professor will not be teaching more than the maximum teaching credits when the proposed section is added to their teaching assignment. Professor is not already teaching a section at the same time as (9) the proposed section. * Oct. to Dec. 1995 solutions by Basbara von Halle in Database Programming and Design A-R01 7472

		Pr	ofe	essor A Instan		-	ime	nt		
CJI №	Department Prefor	Course	Section No.	Course Title No.	Credit	Day	Jime	Bailding	Room	listracion
	MATH MATH	121 845		College Algebra Istro to Prob & Sta			0800-045 1100-123			Sulf W. <u>Zimmer</u>
										•
•• 1995-	96 UNIN cou	rse catal	9							Ja XI.M. Pap 11







Conclusion Natural Language Modeling may be used to analyze any business rule. All business rules may be specified with set theory constraints against elementary sentences. The analyzed facts may be validated by any subject matter expert. The implementation may be tested against the validated requirements. Accountability may be assigned for all aspects of the project. Productivity improves when applications are built according to precise requirements.



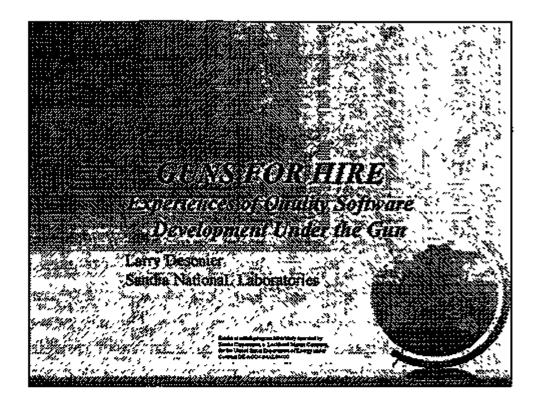
Session A5: Software Quality: Experiences & Year 2000

Chair Cathy Kuhn AS/FM&T

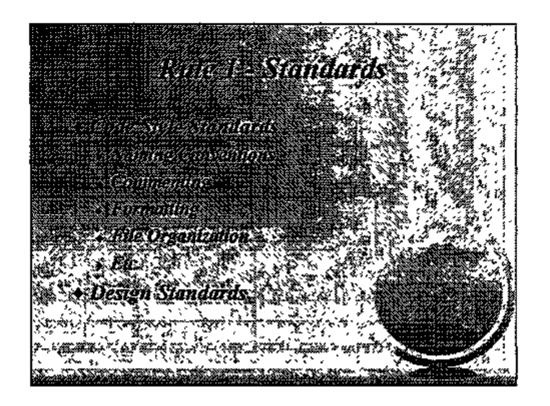
I

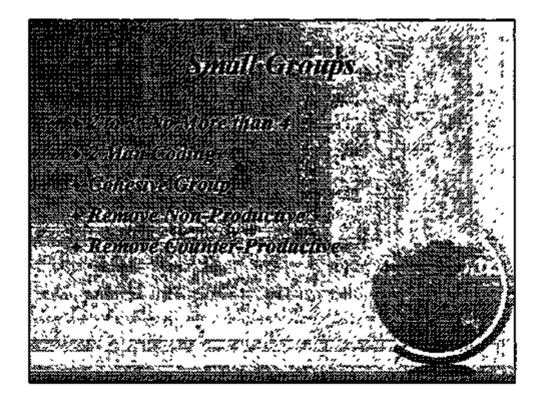
ţ

Session : Paper #	Author(s)	Title
A5:1	Larry Desonier Sandia National Laboratories	Guns for Hire - Experiences of Quality Software Development Under the Gun
A5:2	Bruce Johnston Pantex Plant	The Year 2000 Challenge: A Project Management Perspective
A5:3	Curt Holmes Lockheed Martin Energy Systems	Year 2000 Awareness



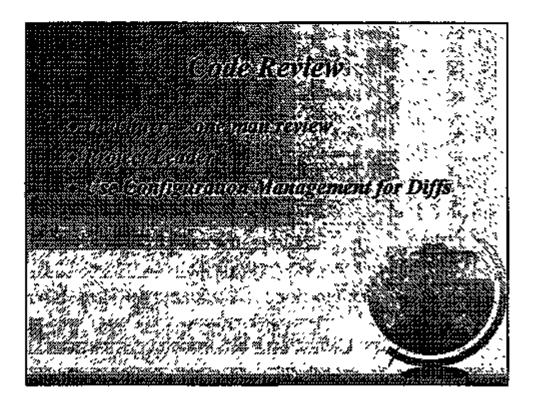


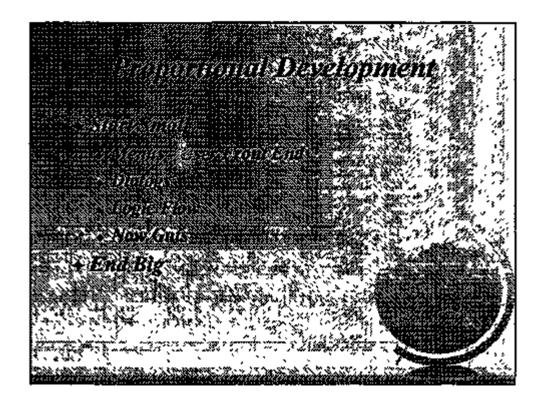


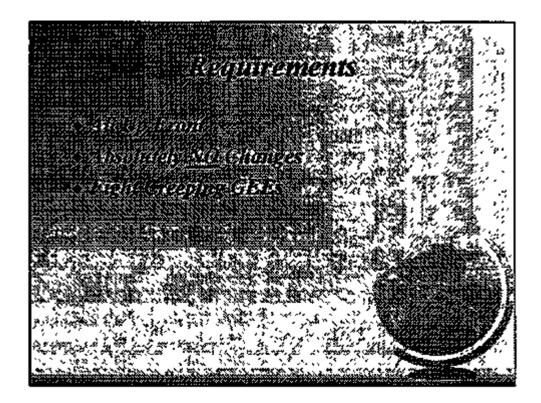


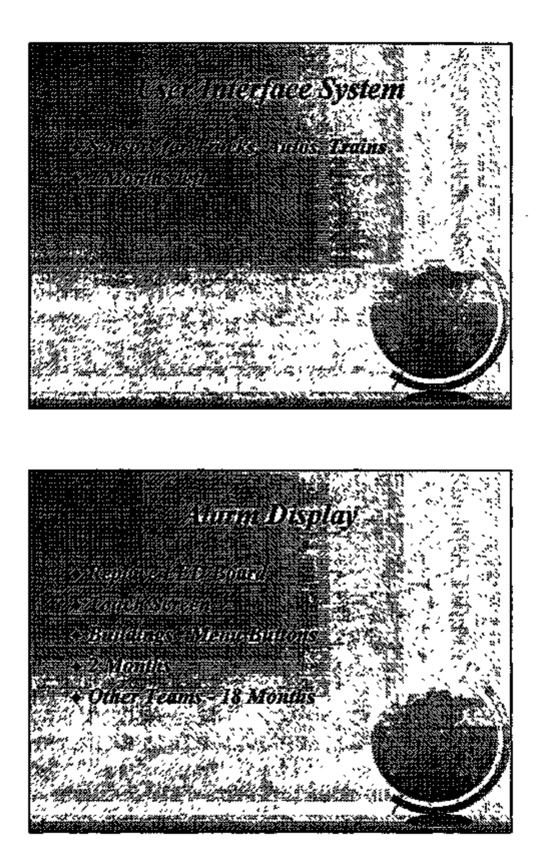
anagement on Management n (normalise) Normalise N 19 12 (42 , **2,813**) ar 2

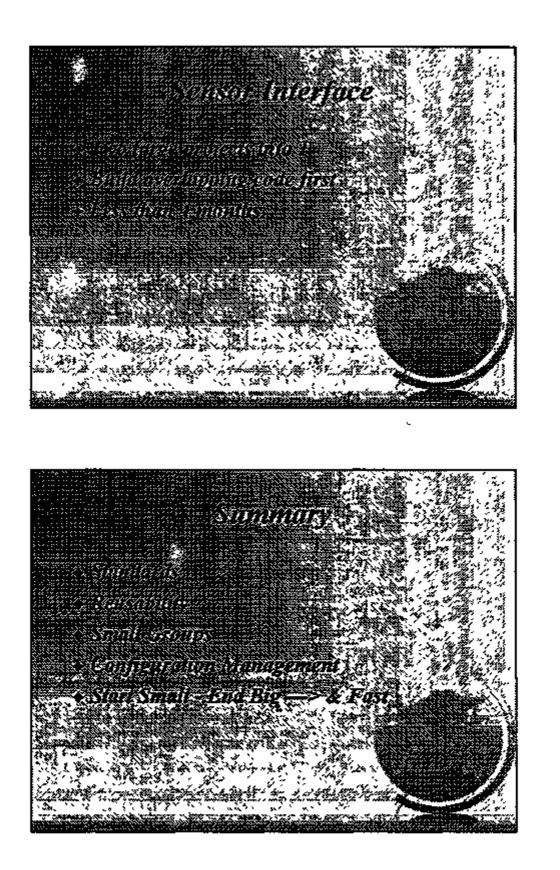
۱,



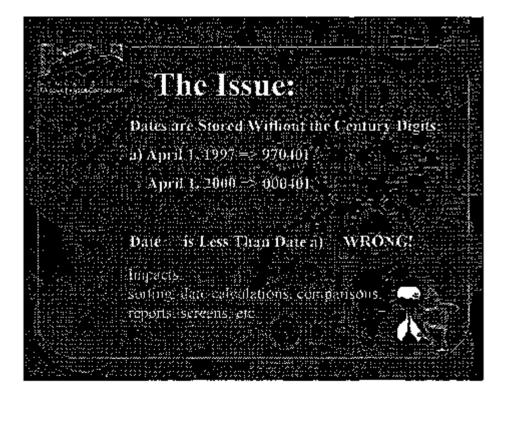


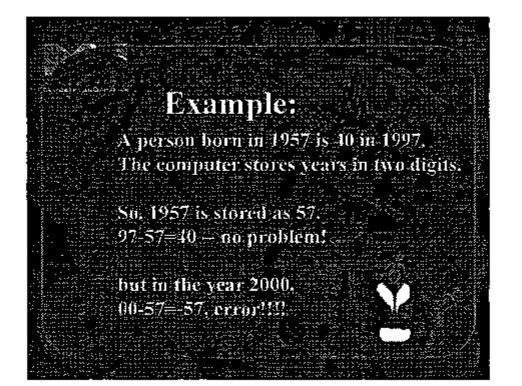


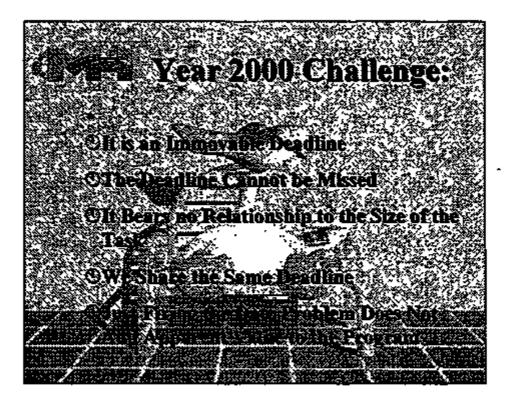




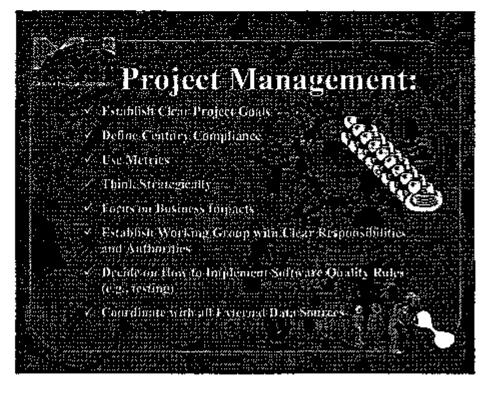


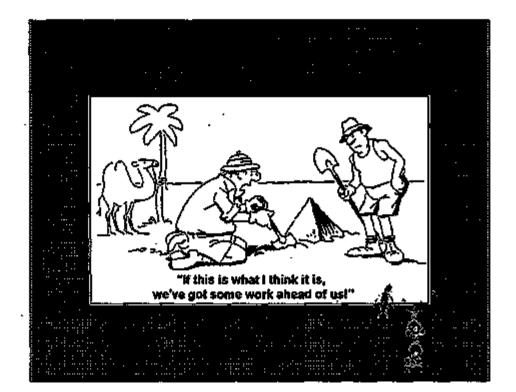






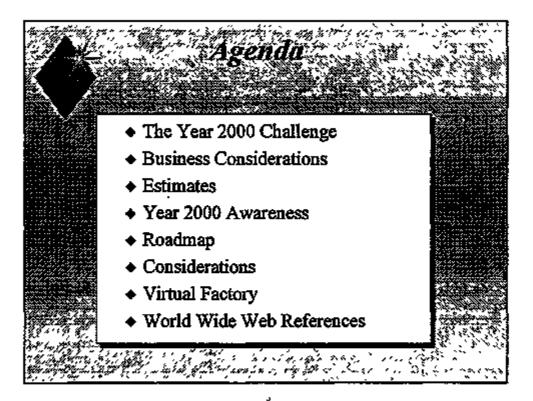


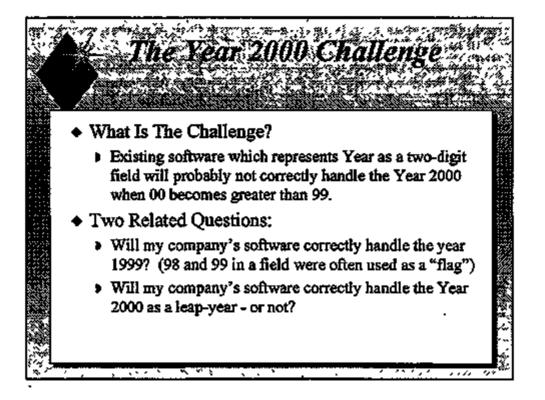


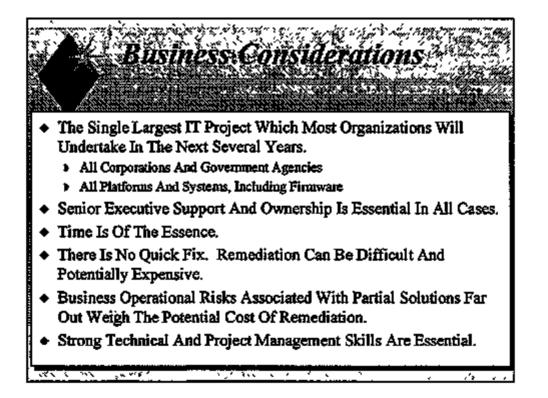


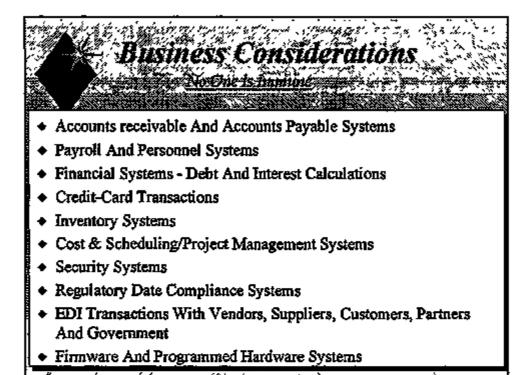
:

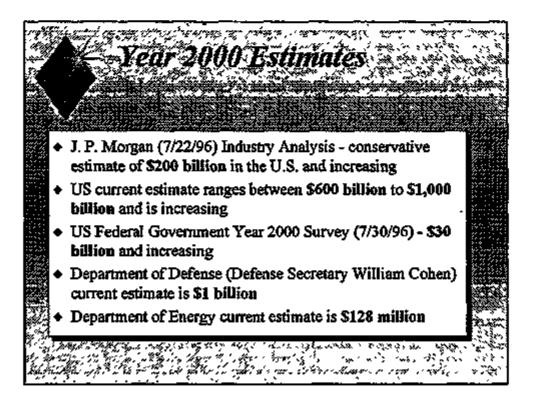


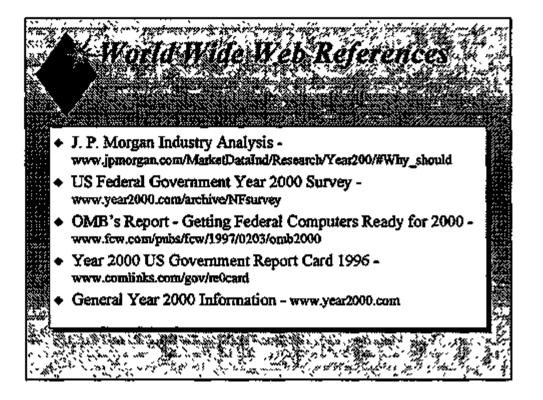


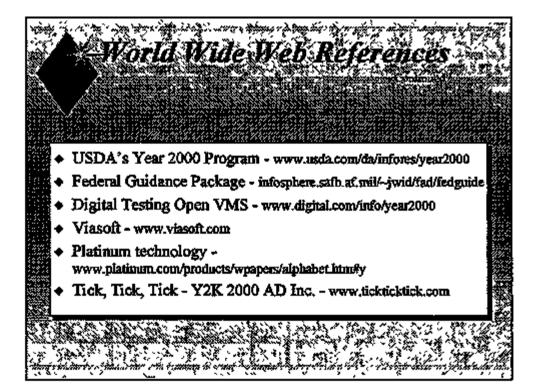


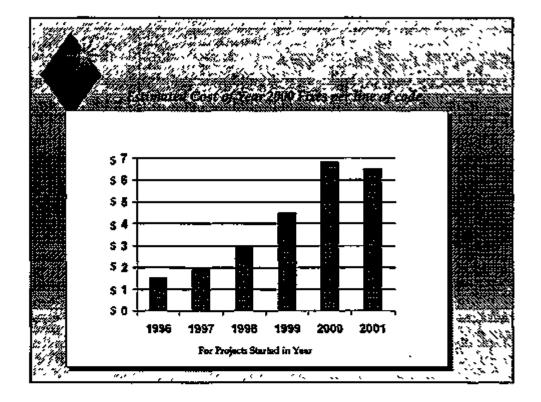


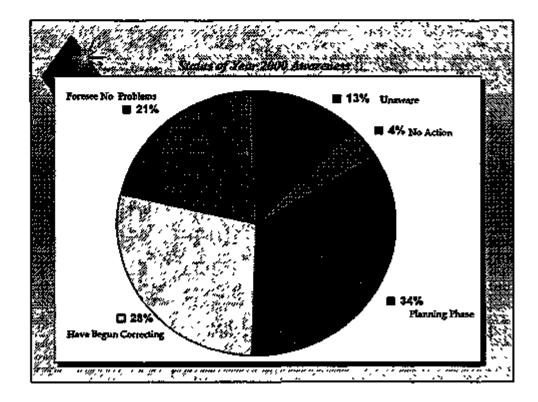


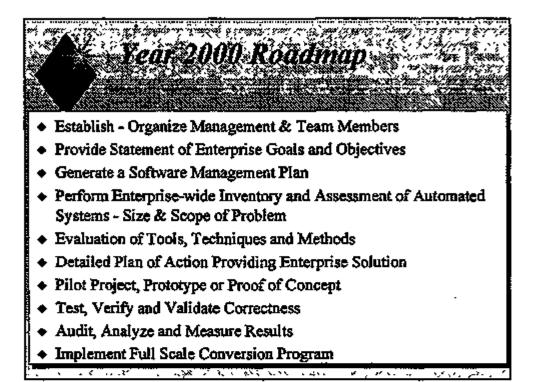


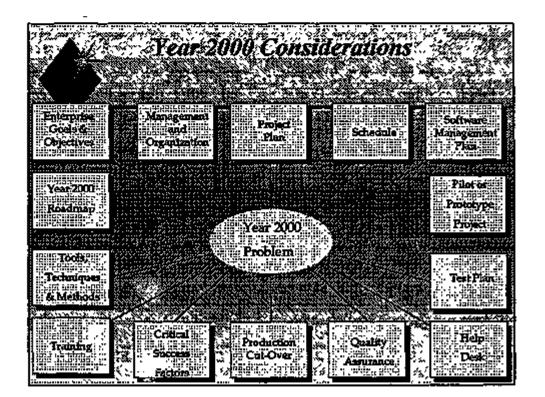


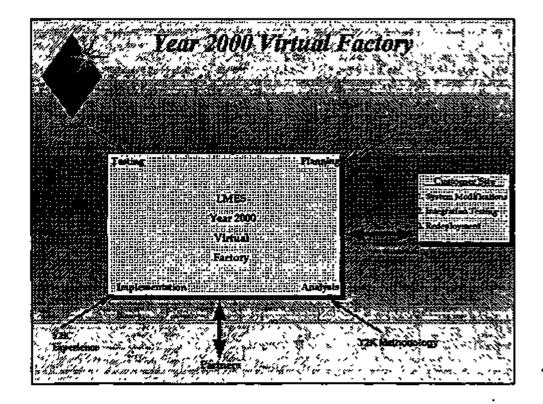


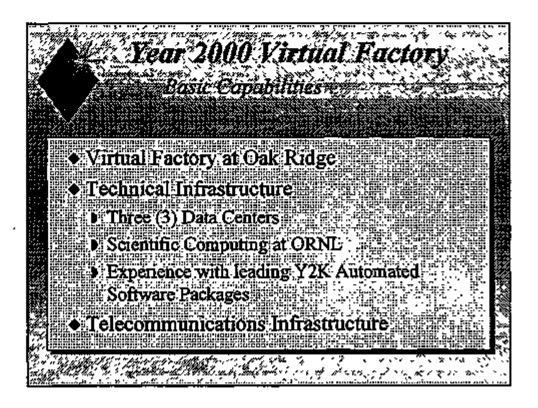


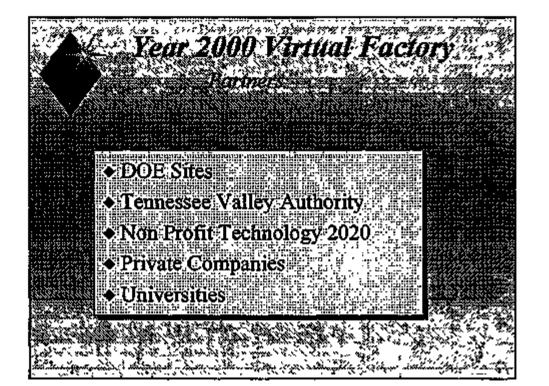


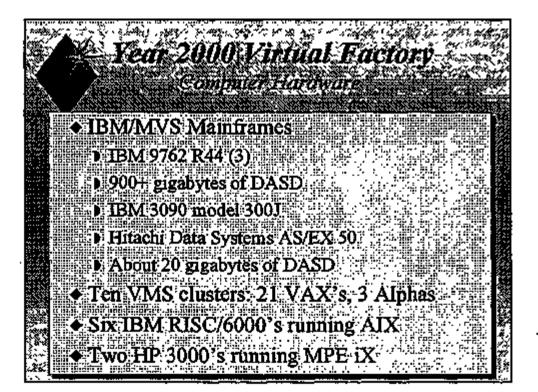


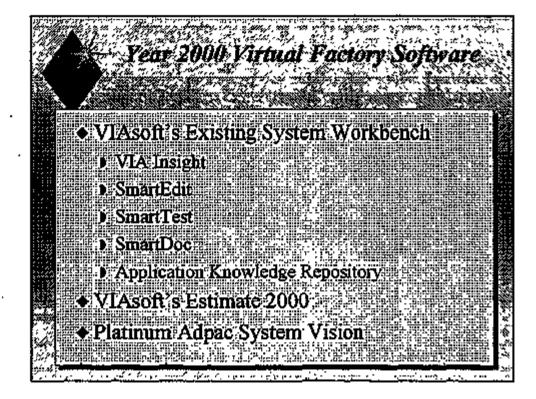


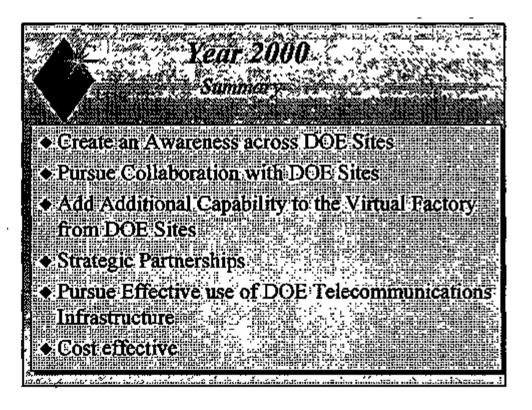










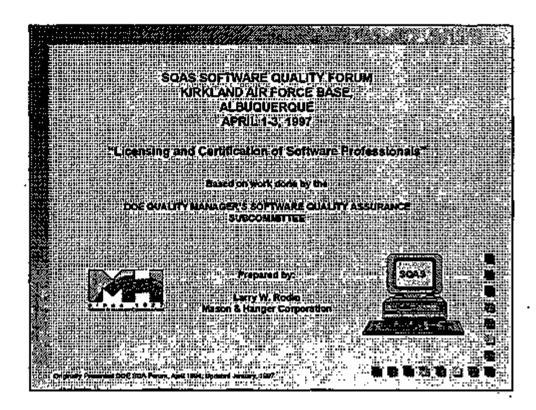




Session B5: Software Standards for Quality Engineering

Chair Patty Trellue Sandia National Laboratories

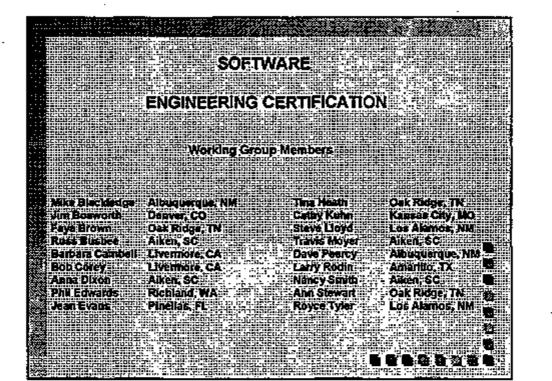
Session : Paper #	Author(s)	Title
B5:1	John Hare AWE UK	ISO and Software Quality Assurance
B5:2	Larry Rodin Pantex Plant	Licensing and Certification of Software Professionals
B5:3	Michael Lackner AS/FM&T	Operational Excellence (Six Sigma) Philosophy: Application to Software Quality Assurance



Licensing and Certification of Software Professionals
 Background for Presentation
 Certification Programs
 Licensing Programs
 Why Become Certified?
 Certification as a Condition of Employment
 Certification Reguirements

Examination Structures

1



SOFTWARE ENGINEERING CERTIFICATION WORKING GROUP

Research softwars-related certification and licensing efforts

Provide (periodic) status reports to the Quality Managers concerning Certification, showing trends from previous reports

Completed Deliverables:

Objectives:

White paper on licensing and certification of software professionals.

Dynamic Resource Notebook on "Software Professionals" certification a and licensing programs: scope (categories/target groups), bodies of knowledge, resource requirements

DEFINITIONS

Cempication —Formal recognition granted by a procession that an individual has demonstrated a proficiency within, and a comprehension of, a specific. Body of Knowledge at a point in time.

License — Permission granted by a government authority to an Individual to engage in a business or occupation or in an activity otherwise unlawful

٩,7

Established Programs Established Programs Institute for Certification of Computer Professional (ICCP) Associate Computer Professional (ACP) Certified Computer Professional (CCP), effective 1/194 Before 1/194, the following designations were affered: Certified Computer Professional (CCP) Certified Data Professional (CCP) Certified Data Professional (CCP) Certified Data Professional (CCP) Certified Computer Professional (CCP) Certified Computer Professional (CCP) Certified Data Professional (CCP) CERTIFIED Data Professional (CCP) CERTIFIED Data Professional

LICENSING OF SOFTWARE ENGINEERS

Gary Ford, Software Engineering Institute (SEI) Technical Staff, presented a paper of the 1993 SET Software Engineering Symposium enflied. The Career State of Certification & Ucerssing of Software Engineers 1. This paper contained accepts on professional Scattaing from these states Perseyfuting, West Wights, and New Versey. NEW UprScr Was the only state Identified at Software States and New Versey. NEW UprScr Was the only state Identified at Software States Software Descenteers Jogialistics (State of New Versey). Assembly Bill 4414 New Versey Software Designeers (Scattaing Sei).

MOTIVATION for LICENSING ENGINEERS

- West Virginia Similar to esteplized bit, hadmor property and to promote the public :

LICENSING ENGINEERS IN OTHER STATES

Members of the SQAG Work Group Died to differmine Software Engineering/Development ICensing afforts in their respective state? California, Colorado, Ronda, Mesouri, New Mexico, Ohio, South Carofina, Tennessee, and Texas. No enidence was found to document Acensing afforts to any of these states

WHY BECOME CERTIFIED? ASOC: • In today's world where quality competition is a

reality, and the need for high-quality software a central concern of many organizations, certification serves as a mark of excellence by demonstrating that the certified individual has the knowledge needed to improve the quality of software. Over 125 organizations have formally recognized ASQC Certification as verification of an individual's possession of this knowledge. Certification is an investment in your career and in the future of your employer.

WHY BECOME CERTIFIED?

ICCP:

Certification is the way to the top of the computing profession. And the prestigious CCP designation Cartifled Computing Professional from ICCP is recognized worldwide by employets and parts as validation of the bolders' Computing Monifestige and Experience. The CCP is the standard which others cover. That is because PCCP, the institute for Gentleation of Computing Professional is administrated by the operation of the institute for and schoology sectors as the most important source of professional cartiflation. Our CCP attantiation demands a high degree of professional competence from those who pass, consequently, the designation is powerful evidence of the high level of attainment of a true Cartifled Computing Professional.

IGGP is the standard to processional certification for 22 indicate and international professional computing accieties - and for numerous individual employers.

Certification is the confidence-building proof that you have not specific requirements and possess high levels of knowledge and skills. And it is easier that ever to become certified, with the introduction or our innovative computer-based festing concept

an bough ecohomic times, certification, adds to your professional oradibility and gives you an advantage in the competitive job market. The recognition shat comes with the CO designation makes ICCP the feducity's leading professional certification organization.

CERTIFICATION AS A CONDITION OF EMPLOYMENT

Equal Employee Opportunity (EEO) laws are detailed regulations published by the federal government which control the employer's use of selection procedures:

If procedures (size) as written tests) affect designated population subgroups, then the exploser must have substantial evidence that the procedure meets a busites necessity.

With paper and peopli tests, adverse effects will normally be assumed unless the employee that avidence to the convery since the results of most tests do differ among population subgroups. Most tests used to education and employment show differences among population subgroups.

An employer has one of two ways to show the procedure or test measure skills about the pob in question.

Differ statistical evidence, usually correlations between test scores and measures of scitual job performance which above that higher scores are anked to higher levels of performance.

2. Show that the content of the exam covers specific job skills which are assemble to the job st question.

<u>+ PRIANAANSENSENDEDAANNEDI</u>

籱

Χ.

邀

OVERVIEW OF THE ICCP REQUIREMENTS FOR ASSOCIATE COMPUTING PROFESSIONAL

- Experience: Any person who has obtained basic knowledge of Information Processing and one of the recognized programming languages may apply for the exam.
- Examination; Pass a two part examination;
 - 1) Core Examination
 - 2) Option of one of eight programming, language
 - Ada, BASIC, C. COBOL: Fortran, Pascal, RPG II,

. • • • •

נוניניים איזייניינייי

28

28.

计多数 计算机

COR Codes: Candidates must subscribe to Code of Ethics,
 Conduct and Good Practice,

OVERVIEW OF THE ICCP REQUIREMENTS FOR CERTIFIED COMPUTING PROFESSIONAL

At months of full-time (of part-time equivalent) protessional experience: A backelor's of graduate depres in 15 or CS or CP Certification may be counted as 24 months experience. A backelor's or graduate depres in a restand field may be counted as 15 months experience. A backelor's or graduate degree to an unretained field may be counted os 12 months appendiate.

of professions limit in Statements from professions collegues stating to

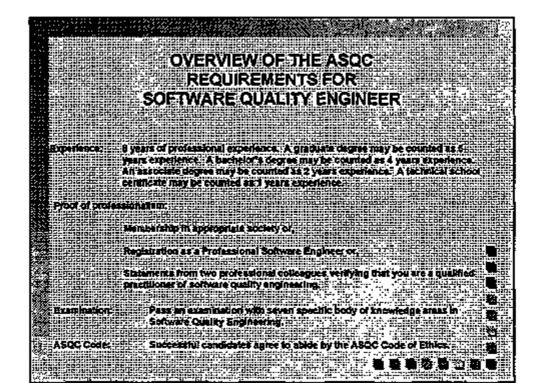
ation: :: Pass & three-part examination;

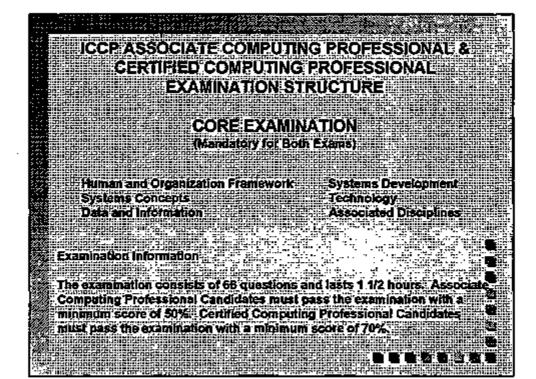
- 3

÷

1) Core Examination 20) Two Speciality Examinations: Management, Procedural/Programming; Systems Development, Business Information Systems, Communications, Office Internation Systems, Systems Security, Software Engineering; Systems Programming, and Data Resource Management

ICCP Codes: ... Candidates must subscribe to Code of Ethics, Conduct and Good Practice





	ICCP ASSOCIATE COMPUTING PROFESSIONAL LANGUAGE EXAMINATION STRUCTURE											
	Choose one langu - C	age examination for Pascal	ACP designation.									
	RPG II	RPG/400	COBAL									
	Fortran	Ada										
	Exemination information		着 - 着 法									
· · ·	The Core Examination con Examination consists of 6 Candidates must pass ear order to receive the ACP of	6 questions. Each Exam th examination with a min	jasts 1 1/2 hours.									

	ICCP CERTIFIED COMPUTING PROFESSIONAL EXAMINATION STRUCTURE											
	(Choose two from following section for CCP designation)											
	Management	Software Engineering										
	Procedural Programming	Communications										
	Systems Development	Office Information Systems										
	Business Information Systems	Systems Programming										
	Systems Security	Data Resource Management 🗮										
	Examination Information											
Y N I		巻 of 110 multiple choice questions each, 会 nours. Candidates must pass both the 義 巻巻を始合ごを始										

- ,'

ŧ

-

ASQC SOFTWARE QUALITY ENGINEER BODY OF KNOWLEDGE General Knowledge, Conduct, л^и Software Metrics, and Ethics (24 Questions) Measurements and Analytical Methods (24 Questions) 1 Software Quality Management (18 Questions) VI. Software Inspection, Software Processes ĨÆ Testing, Verification (24 Questions) and Validation (24 Questions) 11 Software Project Management (16 Questions) Software Audits (16 Questions) Examination Information The Software Quality Engineering exam consists of multiple choice questions. The exam lasts 4 hours, Candidales must pass the exam to 8 be certified.

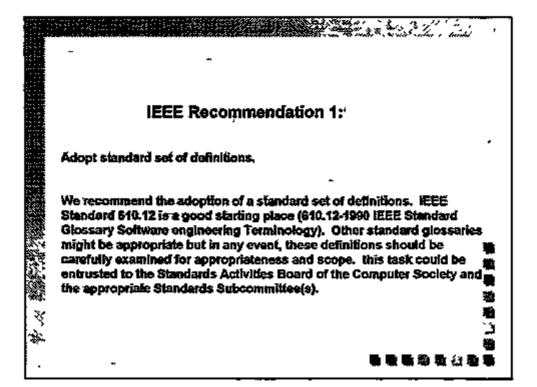


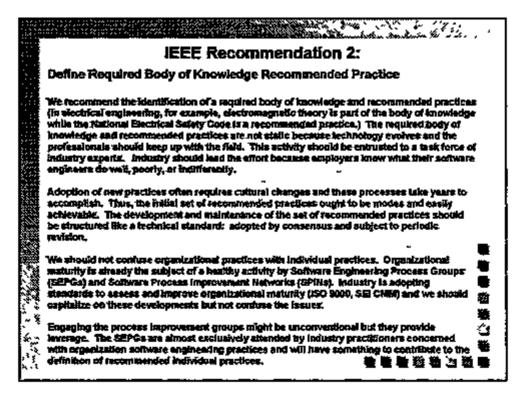
Steering Committee Report: Establishment of Software Engineering as a Profession

Recommendation 1: Adopt Standard Set of Definitions

Recommendation 2: Define Required Body of Knowledge

Recommendation 3: Define Ethical Standards





LEEE Recommendation 3:

Define Ethical Standards

We recommend to study and customize. If necessary, existing codes already adopted by IEEE: ACM, registration boards, and officer relevant organizations: It is not clear that we need something terribly different or specific to software on the ground that the code of ethics of professionals building antennas, processors, or databaset should be different. However, due perhaps to the rapid expansion of the field, software developers sometimes do things that might be considered unethical in other fields (e.g., indiscriminate copying of software in violation of copyrights or licenses.) This back should be charged to the Committee on Public Policy (COPP) of the Computer Society.

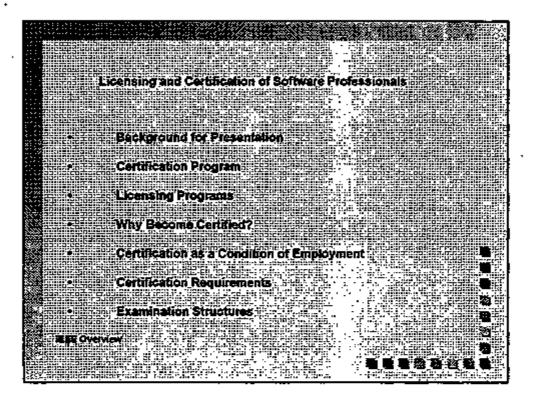
m

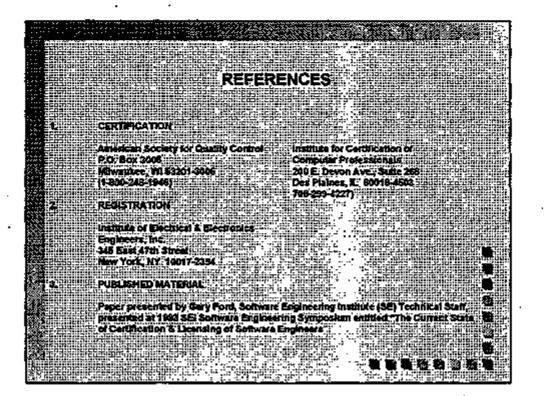
IEEE Recommendation 4:

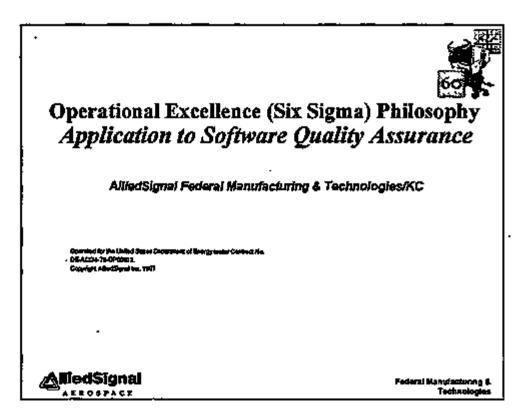
Define Educational Curricula

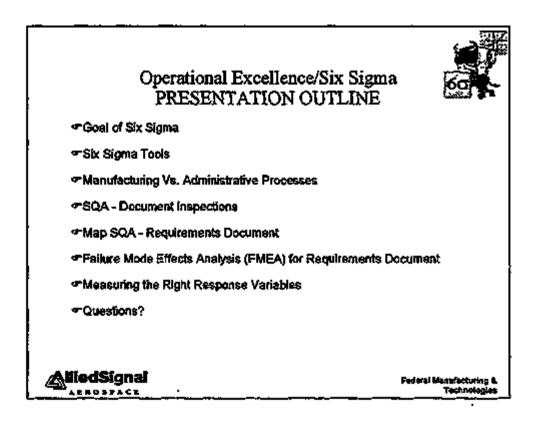
We recommend the definition of ourricula for (a) undergraduate; (b) graduate (MS), and (c) continuing education (for retraining and intgration). This should be charged to an academic task force drawn from educational boards within the SEI, ACM and IEEE Computer Society, and relevant affiliate societies.

There is a debate as to whether Software Engineering is part of Computer Science or vice versa. We should not be distracted by this debate fromt he goal of meeting the needs of industry. The education needed by competent Software engineers could be acquired in different ways. For example, we might identify the need for a foundation on statistics; at a given school, the courses could be offered by Computer Science, Software Engineering, or other departments. The objective is to seek agreement on the curricula that should be faught and not necessarily on which departments beach it:



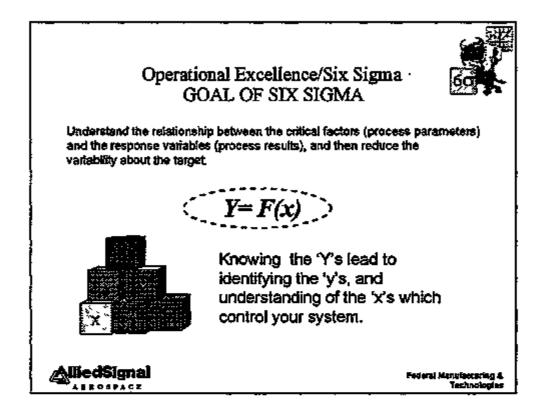


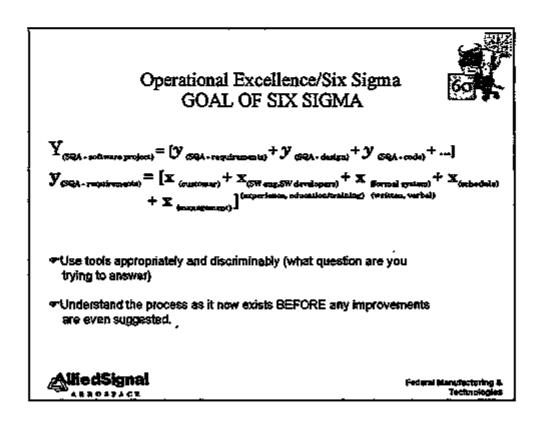


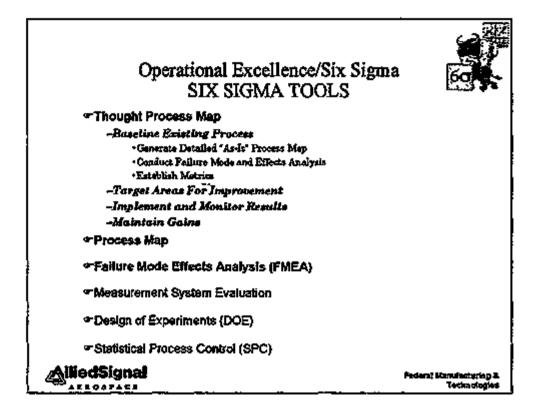


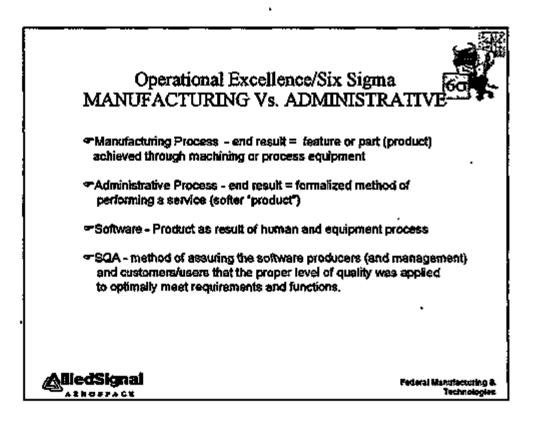
-~',

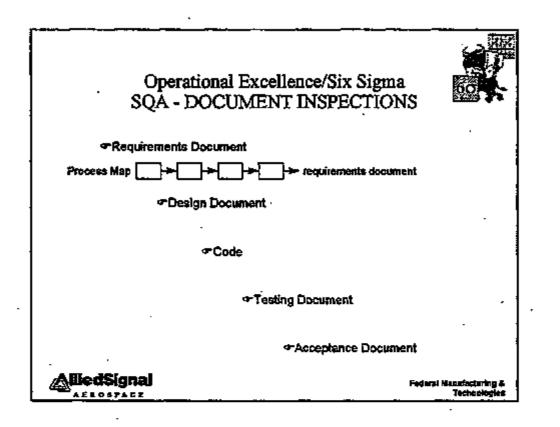
D

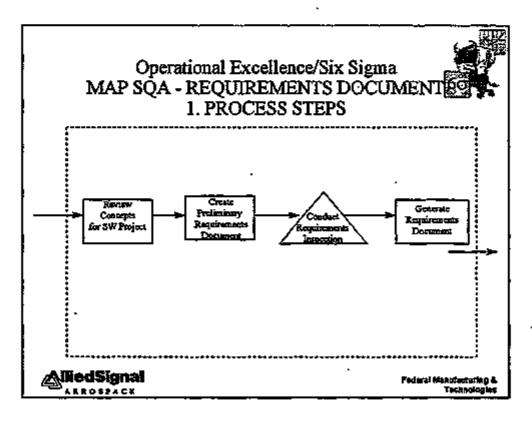


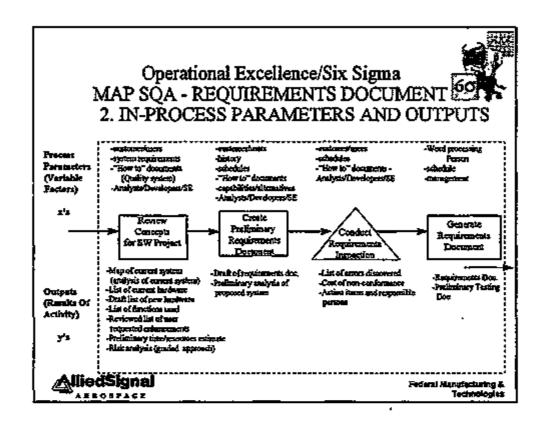


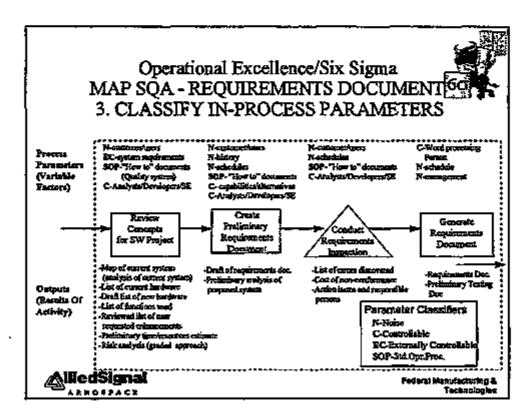






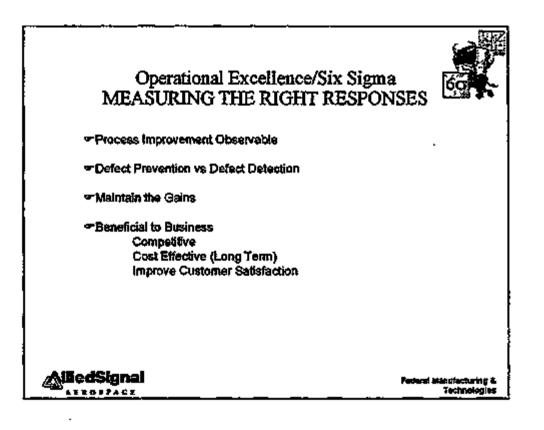


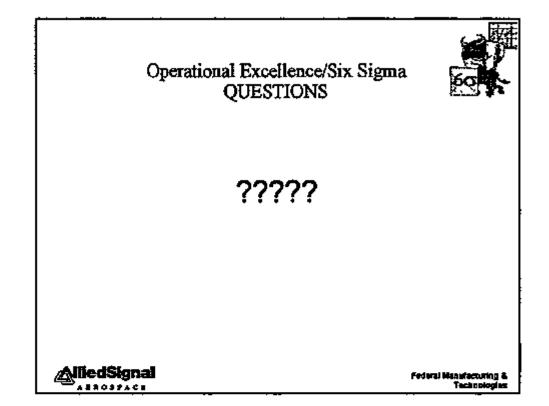




/- -

										Ś	ĉ	2	
	Оре	erationa	al	Excellence	/	Six Sig	n	a		T Sea	3		ľ
FME				EQUIREM						NP	ŝ,	Ş	÷
								-					
					ŀ		•	2		1			!
Talapara	Fallers Made	Referen 200ale		(Canton	ķ		ĥ	F	1.00	Rear			
terne Conspil le Calente Frank	Nara Connei Brian	Dariber/Lete	T,	Last & Splins (In Freedom Teeld) Last by Stilling	ļ		I,				Π	T	
		Deby/mm/mis		Directorian d'andreas and an Thursday	Ī	-1.07 0* Decases	Ī,	Ľ			Π		ļ
	Jugert Tagger Magnage Page	Lake spet	Γ,	Com y 794 as	,		,		Derekter Zaline President				
	Constant Ampleont	lanere barg	Ī,	Lais a' Coisean i an Eir-main	ļ		ļ,		1		Π	Ì	Ī
	Leig af Page pageages Proppie Dafasian	Delegistași ente	[,	Paur pro pour ley Solito aire Engraine Lymp (Daaring ay	Ļ		ļ	æ					1
			L		Ŀ		ł			{			
		<u> </u>	┡		Ļ		Ļ		ļ	<u> </u>	Ц	+	Ł
			⊢		ł	· ·	ł	┝		 .	Н	╈	╋
			1.		t		ŀ		· · · · ·		Н	T.	t
					ſ		ſ	-		<u>ا</u>	П	T	Ĺ
			┡		Ļ	· · · · ·	Ľ			<u> </u>	Н	4	╇
	dSignal	I	I	<u>. </u>	ł,	1		Ļ	Federal	i. Voorte Tech			





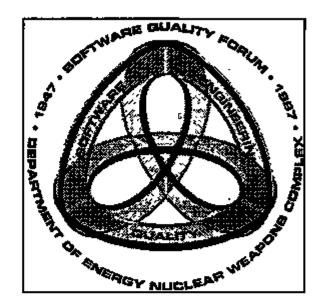
t

1 Î

ł

, t•

.



Wrapup and Awards

Best Tutorial Award Best Presentation Award