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## RANDOMIZED PUSH-OUT MECHANISMS IN PRIORITY QUEUEING AND THEIR PROBABILITY CHARACTERISTICS

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ABSTRACT. The non-preemptive priority queueing with a finite buffer is considered. A randomized push-out buffer management mechanism that allows to control very efficiently the loss probability of priority packets is introduced. The packet loss probabilities for priority and non-priority traffic are derived with the use of the generating function approach. For the standard non-randomized push-out scheme, the explicit analytic expressions are obtained. A procedure for the numerical calculation of mean queues is also proposed.

### 1. Introduction

Priority queueing disciplines have a number of important applications in computer networks, for example, in the Differentiated Services architecture for the Internet [7].

Consider the non-preemptive priority queueing system with two classes of packets. Class 1 packets have priority over class 2 packets. The packets of class 1 (2) arrive into the buffer according to the Poisson process with rate  $\lambda_1$  ( $\lambda_2$ ), respectively. The service time has the exponential distribution with the same rate  $\mu$  for each class. The service times are independent of the arrival processes. The buffer has a finite size  $N$  and it is shared by both types of

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customers. If the buffer is full, a new coming customer of class 1 can push out of the buffer a customer of class 2 with the probability  $\alpha$ . Note that if  $\alpha = 1$  we retrieve the standard non-randomized push-out mechanism.

The infinite buffer priority queueing has been thoroughly studied in [4, 8, 9]. The case of finite buffer priority queueing received considerably less attention. The M/M/C/K type finite buffer non-preemptive priority queueing with non-randomized push-out mechanism is analyzed by Kapadia *et al* [5, 6]. Bondi [1] considers the M/M/1/K type preemptive and non-preemptive priority queueing with the following buffer management schemes: complete partitioning, complete sharing and sharing with minimum allocation. Wagner and Krieger [10] analyze the M/M/1/K type non-preemptive priority queueing with the complete sharing buffer management scheme and with the class-dependent service rates. In [2] Cheng and Akyildiz consider the priority queueing with general service time distributions and a general service discipline function.

Most of the above works use recursive relations to solve steady state Kolmogorov equations. We use a generating function approach, which only requires the solution of a linear system of  $N$  equations. As far as we know, the randomized push-out mechanism is analyzed for the first time.

## 2. The generating functions

Denote by  $p(i, n)$  the stationary probability of the event that there are  $n$  packets in the queue including  $i$  packets of class 1. We also use  $p_0$  for the stationary probability of the event that there are no packets in the system. These probabilities satisfy the following stationary Kolmogorov equations:

$$(\lambda_1 + \lambda_2)p_0 = \mu p(0, 0);$$

- $n = 0$

$$(\lambda_1 + \lambda_2 + \mu)p(0, 0) = \mu p(1, 1) + \mu p(0, 1) + (\lambda_1 + \lambda_2)p_0;$$

- $0 < n < N$

$$\begin{aligned} (\lambda_1 + \lambda_2 + \mu)p(0, n) &= \mu p(1, n + 1) + \mu p(0, n + 1) + \lambda_2 p(0, n - 1), \\ (\lambda_1 + \lambda_2 + \mu)p(i, n) &= \mu p(i + 1, n + 1) + \lambda_1 p(i - 1, n - 1) + \lambda_2 p(i, n - 1), \\ (\lambda_1 + \lambda_2 + \mu)p(i, n) &= \mu p(n + 1, n + 1) + \lambda_1 p(n - 1, n - 1); \end{aligned}$$

- $n = N$

$$\begin{aligned}
 (\alpha\lambda_1 + \mu)p(0, N) &= \lambda_2 p(0, N-1), \\
 (\alpha\lambda_1 + \mu)p(i, N) &= \lambda_1 p(i-1, N-1) + \lambda_2 p(i, N-1) + \alpha\lambda_1 p(i-1, N), \\
 \mu p(N, N) &= \lambda_1 p(N-1, N-1) + \alpha\lambda_1 p(N-1, N).
 \end{aligned}$$

Next we introduce the generating function for  $p(i, n)$  by the index  $i$

$$F_n(x) = \sum_{i=0}^n p(i, n)x^i.$$

Using the above given Kolmogorov equations, we obtain the following relations for the generating functions  $F_n(x)$ ,  $n = 0, 1, \dots, N$ :

- $n = 0$

$$(\lambda_1 + \lambda_2 + \mu)F_0(x) = \frac{\mu}{x} [F_1(x) - p(0, 1)] + \mu p(0, 1) + (\lambda_1 + \lambda_2)p_0,$$

- $0 < n < N$

$$(\lambda_1 + \lambda_2 + \mu)F_n(x) = \frac{\mu}{x} [F_{n+1}(x) - p(0, n+1)] + \mu p(0, n+1) + (\lambda_1 x + \lambda_2)F_{n-1}(x).$$

In particular, we get the following boundary condition

- $n = N$

$$\begin{aligned}
 (1) \quad (\alpha\lambda_1 + \mu)F_N(x) - \alpha\lambda_1 p(N, N)x^N &= (\lambda_1 x + \lambda_2)F_{N-1}(x) \\
 &\quad + \alpha\lambda_1 x F_N(x) - \alpha\lambda_1 x^{N+1} p(N, N).
 \end{aligned}$$

Now we introduce the generating function for  $F_n(x)$  by the index  $n$

$$\Phi(x, y) = \sum_{n=0}^{N-1} F_n(x)y^n.$$

The generating function  $\Phi(x, y)$  satisfies equation (2) given in Lemma 1 below.

**Lemma 1.** *The generating function  $\Phi(x, y)$  satisfies the following equation*

$$\begin{aligned}
 (2) \quad [(\rho + 1)xy - xy^2(\rho_1 x + \rho_2) - 1]\Phi(x, y) &= -y^{N+1}x(\rho_1 x + \rho_2)F_{N-1}(x) \\
 &\quad + y^N F_N(x) + y(x-1)A(y) + (xy-1)\rho p_0,
 \end{aligned}$$

where  $\rho_i = \lambda_i/\mu$ ,  $\rho = \rho_1 + \rho_2$  and  $A(y) = \sum_{n=0}^{N-1} p(0, n+1)y^n$ .

The generating function  $\Phi(x, y)$  is determined by the next result.

**Theorem 1.** *The generating function  $\Phi(x, y)$  is given by*

$$\begin{aligned} \Phi(x, y) = & \\ & \frac{[1 - xy + \alpha\rho_1xy(x - 1)]y^N V_{N-1}(x) + y(x - 1)A(y)}{(\rho + 1)xy - xy^2(\rho_1x + \rho_2) - 1} \\ & + \frac{[1 - xy]x^N y^N p(N, N) + \rho[xy - 1]p_0}{(\rho + 1)xy - xy^2(\rho_1x + \rho_2) - 1}, \end{aligned}$$

where

$$\begin{aligned} V_{N-1}(x) &= \sum_{k=0}^{N-1} x^k p(k, N), \\ A(y) &= -\alpha\rho y^{N-1} p(0, N) \\ &+ \sum_{k=1}^{N-1} \left[ \rho_2 y^{N-k} \frac{U_{k-1}(t)}{\rho_1^{(k+1)/2}} - \alpha\rho y^{N-k-1} \frac{U_k(t)}{\rho_1^{k/2}} \right. \\ &\left. + \alpha y^{N-k-1} \frac{U_{k-1}(t)}{\rho_1^{(k-1)/2}} \right] p(k, N) + \rho_2 \frac{U_{N-1}(t)}{\rho_1^{(N+1)/2}} p(N, N) \end{aligned}$$

with  $t = (\rho + 1 - \rho_2 y) / (2\rho_1^{1/2})$  and where probabilities  $p(k, N)$ ,  $k = 0, \dots, N$  can be obtained as a solution to the following system of linear equations

- $s = 0$

$$\begin{aligned} & \alpha\rho_1 C_{N-1}^1(t_0) p(N - 1, N) \\ & + \left[ \rho C_{N-1}^1(t_0) - \rho_1^{1/2} C_N^1(t_0) \right] p(N, N) + \rho\rho_1^{(N+1)/2} p_0 = 0, \end{aligned}$$

- $0 < s < N$

$$\begin{aligned} & \sum_{k=0}^{s-1} \left[ \rho \frac{C_{N-s-1}^{s-k}(t_0) \rho_1^{k+1}}{(-\rho_2)^{k+1}} - \rho_1^{3/2} (1 + \alpha\rho) \frac{C_{N-s}^{s-k}(t_0) \rho_1^k}{(-\rho_2)^{k+1}} \right. \\ & \left. + \rho_1 \alpha \frac{C_{N-s-1}^{s-k+1}(t_0) \rho_1^k}{(-\rho_2)^k} \right] p(N - 1 - k, N) + \alpha\rho_1^{s+1} \frac{C_{N-s-1}^1(t_0)}{(-\rho_2)^s} p(N - 1 - s, N) \\ & + \left[ \rho C_{N-s-1}^{s+1}(t_0) - \rho_1^{3/2} C_{N-s}^{s+1}(t_0) \right] p(N, N) = 0, \end{aligned}$$

- $s = N$

$$-\rho_1^{3/2}(1+\alpha\rho) \sum_{k=0}^{N-1} \frac{C_0^{N-k}(t_0)\rho_1^k}{(-\rho_2)^{k+1}} p(N-1-k, N) - \rho_1^{1/2} C_0^{N+1}(t_0) p(N, N) = 0$$

with  $U_n(x)$  and  $C_n^\nu(x)$  denoting the Chebyshev polynomials of the second kind and the Gegenbauer polynomials [3], respectively, and

$$p_0 = (1 - \rho)/(1 - \rho^{N+2}), \quad t_0 = (\rho + 1)/(2\rho_1^{1/2}).$$

The proof is carried out into Section 4.

### 3. The loss probabilities

Once we know the value of  $p(N, N)$ , we can derive the loss probabilities of class 1 and class 2 packets.

**Theorem 2.** *The loss probabilities of class 1 and class 2 packets are given by the following formulae*

$$(3) \quad P_{loss}^{(1)} = p(N, N) + (1 - \alpha)[P_N - p(N, N)],$$

$$(4) \quad P_{loss}^{(2)} = P_N + \alpha \frac{\rho_1}{\rho_2} [P_N - p(N, N)],$$

where

$$P_N = \frac{1 - \rho}{1 - \rho^{N+2}} \rho^{N+1}.$$

**Proof:** A priority packet can be lost either when the whole buffer is filled only with priority packets or when there are some packets of class 2 but with probability  $1 - \alpha$  the push-out mechanism is not enabled. The probability of the first event is  $p(N, N)$  and the probability of the second event is  $\sum_{k=0}^{N-1} p(k, N) = P_N - p(N, N)$ . Thus, we obtain formula (3).

The stream of lost packets of class 2 consists of the stream of packets with rate  $\lambda_2 P_N$  lost when the buffer is full and the stream of packets with rate  $\alpha \lambda_1 (P_N - p(N, N))$  pushed out by packets of class 1. Since the system is ergodic, we obtain formula (4).

Note that if  $\alpha = 0$  (no push-out), the loss probabilities of two classes coincide and are equal to  $P_N$ . Furthermore, due to the fact that the service time distribution is the same for the two classes, the expressions for  $p_0$ ,  $F_N(1)$  and  $\Phi(1, 1)$  could be obtained immediately by elementary considerations.

In the particular case of the non-randomized push-out mechanism, that is, when  $\alpha = 1$ , we can calculate the loss probabilities explicitly.

**Theorem 3.** *The loss probabilities of class 1 and class 2 packets in the case of non-randomized push-out mechanism are given by*

$$(5) \quad P_{loss}^{(1)} = \rho \rho_1^N \frac{(1 - \rho_1)(1 - \rho^{N+1})}{(1 - \rho_1^{N+1})(1 - \rho^{N+2})},$$

$$(6) \quad P_{loss}^{(2)} = P_N + \frac{\rho_1}{\rho_2} [P_N - P_{loss}^{(1)}].$$

In the case of non-randomized push-out mechanism ( $\alpha = 1$ ), the equation for the generating function (2) takes the form

$$(7) \quad [(\rho + 1)xy - xy^2(\rho_1x + \rho_2) - 1]\Phi(x, y) = y^N[1 - xy + \rho_1x(x - 1)y]F_N(x) \\ + y(x - 1)A(y) + \rho_1(1 - x)x^{N+1}y^{N+1}p(N, N) + (xy - 1)\rho p_0.$$

Setting  $x = 1$  in (7), and then reducing it by the term  $(y - 1)$ , we get

$$(1 - \rho y)\Phi(1, y) = \rho p_0 - y^N F_N(1).$$

Then in the above equation we take subsequently  $y = 1$  and  $y = 1/\rho$  to obtain

$$(8) \quad (1 - \rho)\Phi(1, 1) = \rho p_0 - F_N(1)$$

and

$$(9) \quad 0 = \rho p_0 - \frac{1}{\rho^N} F_N(1).$$

Solving equations (8) and (9) together with the normalization condition

$$\Phi(1, 1) + p_0 + F_N(1) = 1,$$

we obtain the following expressions for  $p_0$ ,  $F_N(1)$  and  $\Phi(1, 1)$ :

$$p_0 = \frac{1 - \rho}{1 - \rho^{N+2}}, \quad F_N(1) = \frac{1 - \rho}{1 - \rho^{N+2}} \rho^{N+1}, \quad \Phi(1, 1) = \frac{1 - \rho^{N+1}}{1 - \rho^{N+2}} \rho.$$

Next we take  $y = 1$  in equation (7) and then reduce it by the term  $(x - 1)$

$$(1 - \rho_1 x)\Phi(x, 1) = -(1 - \rho_1 x)F_N(x) + A(1) - \rho_1 x^{N+1}p(N, N) + \rho p_0.$$

We now set subsequently  $x = 1$  and  $x = 1/\rho_1$  in the above equation. This results in the following two equations:

$$(10) \quad (1 - \rho_1)\Phi(1, 1) = -(1 - \rho_1)F_N(1) + A(1) - \rho_1 p(N, N) + \rho p_0,$$

$$(11) \quad 0 = A(1) - \frac{1}{\rho_1^N}p(N, N) + \rho p_0.$$

Solving equations (10) and (11), we obtain

$$p(N, N) = \frac{(1 - \rho_1)(1 - \rho^{N+1})}{(1 - \rho_1^{N+1})(1 - \rho^{N+2})}\rho\rho_1^N.$$

The loss probability of class 1 packets  $P_{loss}^{(1)}$  is given by  $p(N, N)$ . Then, we note that the stream of lost packets of class 2 consists of the stream of packets with rate  $\lambda_2 F_N(1)$  lost when the buffer is full and the stream of packets with rate  $\lambda_1(F_N(1) - p(N, N))$  pushed out by packets of class 1. Hence, using the ergodicity property of the system, we obtain formula (6) for  $P_{loss}^{(2)}$ .

#### 4. Proof of Theorem 1

By substituting boundary condition (1) into equation (2) for the generating function  $\Phi(x, y)$ , we get

$$(12) \quad [(\rho + 1)xy - xy^2(\rho_1 x + \rho_2) - 1]\Phi(x, y) = [1 - xy + \alpha\rho_1 xy(x - 1)]y^N V_{N-1}(x) \\ + [1 - xy]x^N y^N p(N, N) + y(x - 1)A(y) + \rho[xy - 1]p_0,$$

where  $V_{N-1}(x) = \sum_{i=0}^{N-1} x^i p(i, N)$ , and hence the expression for  $\Phi(x, y)$ .

Next, we set  $z := xy$  and rewrite equation (12) as follows:

$$[(\rho_1 + \rho_2 + 1)z - \rho_1 z^2 - \rho_2 yz - 1]\Phi\left(\frac{z}{y}, y\right) = [(1 - z)y + \rho_1 \alpha(z - y)z]y^{N-1} V_{N-1}\left(\frac{z}{y}\right) \\ + (z - y)A(y) + (1 - z)z^N p(N, N) + \rho(z - 1)p_0.$$

Let us now consider the analyticity condition for the generating function  $\Phi(z/y, y)$ . Namely, the following two conditions have to be satisfied simultaneously

$$(\rho_1 + \rho_2 + 1)z - \rho_1 z^2 - \rho_2 yz - 1 = 0,$$



$$\begin{aligned} & [(1-z)y + \rho_1\alpha(z-y)z]y^{N-1}V_{N-1}\left(\frac{z}{y}\right) \\ & + (z-y)A(y) + (1-z)z^N P(N, N) + \rho(z-1)p_0 = 0. \end{aligned}$$

The first condition can be rewritten as

$$\rho_2(y-z)z = (1-z)(\rho z - 1),$$

which gives

$$y - z = \frac{(1-z)(\rho z - 1)}{\rho_2 z}.$$

We substitute the above expression for  $y - z$  into the first two terms of the second analyticity condition and then reduce it by  $1 - z$ , to get

$$\begin{aligned} (13) \quad & \left(y - \frac{\rho_1}{\rho_2}\alpha(\rho z - 1)\right)y^{N-1}V_{N-1}\left(\frac{z}{y}\right) \\ & - \frac{\rho z - 1}{\rho_2 z}A(y) + z^N p(N, N) - (\rho_1 + \rho_2)p_0 = 0. \end{aligned}$$

Next we denote by  $a$  and  $b$  the roots of the following quadratic equation with respect to the variable  $z$

$$(\rho_1 + \rho_2 + 1)z - \rho_1 z^2 - \rho_2 y z - 1 = 0.$$

Now we substitute subsequently the roots  $a$  and  $b$  into (13), which allows us to eliminate  $A(y)$

$$\begin{aligned} & \frac{\rho b - 1}{b} \left(y - \frac{\rho_1}{\rho_2}\alpha(\rho a - 1)\right)y^{N-1}V_{N-1}\left(\frac{a}{y}\right) \\ & - \frac{\rho a - 1}{a} \left(y - \frac{\rho_1}{\rho_2}\alpha(\rho b - 1)\right)y^{N-1}V_{N-1}\left(\frac{b}{y}\right) + \left(\frac{\rho b - 1}{b}a^N - \frac{\rho a - 1}{a}b^N\right)p(N, N) \\ & - \rho \left(\frac{\rho b - 1}{b} - \frac{\rho a - 1}{a}\right)p_0 = 0. \end{aligned}$$

Taking into account the properties of roots of the quadratic equation

$$ab = 1/\rho_1, \quad (\rho a - 1)(\rho b - 1) = \frac{\rho_2}{\rho_1}(\rho y - 1),$$

we have

$$\left((\rho - \rho_1 a)y - q(\rho y - 1)\rho_1 a\right)y^{N-1}V_{N-1}\left(\frac{a}{y}\right) - \left((\rho - \rho_1 b)y - q(\rho y - 1)\rho_1 b\right)y^{N-1}V_{N-1}\left(\frac{a}{y}\right)$$

$$\begin{aligned}
 & +(\rho(a^N - b^N) - \rho_1(a^{N+1} - b^{N+1}))p(N, N) + \rho\rho_1(a - b)p_0 = 0, \\
 & \rho y^N \left( V_{N-1} \left( \frac{a}{y} \right) - V_{N-1} \left( \frac{b}{y} \right) \right) - \rho_1(y + q(\rho y - 1)) \left( a V_{N-1} \left( \frac{a}{y} \right) - \right. \\
 & \left. b V_{N-1} \left( \frac{b}{y} \right) \right) y^{N-1} + (\rho(a^N - b^N) - \rho_1(a^{N+1} - b^{N+1}))p(N, N) + \rho\rho_1(a - b)p_0 = 0, \\
 & \rho y \sum_{i=1}^{N-1} v_i(a^i - b^i)y^{N-1-i} - \rho_1(y + q(\rho y - 1)) \sum_{i=0}^{N-1} v_i(a^{i+1} - b^{i+1})y^{N-1-i} \\
 (14) \quad & +(\rho(a^N - b^N) - \rho_1(a^{N+1} - b^{N+1}))p(N, N) + \rho\rho_1(a - b)p_0 = 0.
 \end{aligned}$$

By denoting  $\cos \varphi = (\rho + 1 - \rho_2 y)/(2\rho_1^{1/2})$ , the roots  $a$  and  $b$  can be written in the form

$$a = \frac{\exp(i\varphi)}{\rho_1^{1/2}}, \quad b = \frac{\exp(-i\varphi)}{\rho_1^{1/2}}.$$

Then equation (14) can be rewritten as

$$\begin{aligned}
 & \rho y \sum_{i=1}^{N-1} v_i U_{i-1}(t) \frac{y^{N-1-i}}{\rho_1^{i/2}} - \rho_1(y + q(\rho y - 1)) \sum_{i=0}^{N-1} v_i U_i(t) \frac{y^{N-1-i}}{\rho_1^{(i+1)/2}} \\
 (15) \quad & + \left( \rho U_{N-1}(t) \frac{1}{\rho_1^{(N)/2}} - \rho_1 U_N(t) \frac{1}{\rho_1^{(N+1)/2}} \right) p(N, N) + \rho\rho_1^{1/2} p_0 = 0,
 \end{aligned}$$

where  $t := \cos \varphi = (\rho + 1 - \rho_2 y)/(2\rho_1^{1/2})$  and  $U_s(t)$  are the Chebyshev polynomials of the second kind [3]

$$U_s(\cos \varphi) = \frac{\sin(s+1)\varphi}{\sin \varphi}.$$

The Taylor series for the function  $U_s(t)$  with respect to  $y$ , being actually a polynomial in this case, has the following form

$$U_s(t(y)) = \sum_{s=0}^s \frac{U_s^{(i)}(t_0)}{i!} (-1)^i \frac{\rho_2^i y^i}{2^i \rho_1^{i/2}}$$

with  $t_0 = (\rho + 1)/(2\rho_1^{1/2})$ . By changing the order of summation in the expressions

$$\begin{aligned}
 \sum_{i=1}^{N-1} v_i U_{i-1}(t) \frac{y^{N-1-i}}{\rho_1^{i/2}} &= \sum_{l=0}^{N-2} y^l \sum_{k=0}^l v_{N-1-k} \frac{U_{N-k-2}^{(l-k)}(t_0) (-\rho_2)^{l-k}}{(l-k)! 2^{l-k} \rho_1^{(N-1-2k+l)/2}}, \\
 \sum_{i=0}^{N-1} v_i U_i(t) \frac{y^{N-1-i}}{\rho_1^{(i+1)/2}} &= \sum_{l=0}^{N-1} y^l \sum_{k=0}^l v_{N-1-k} \frac{U_{N-k-1}^{(l-k)}(t_0) (-\rho_2)^{l-k}}{(l-k)! 2^{l-k} \rho_1^{(N-2k+l)/2}},
 \end{aligned}$$

we rewrite equation (15) as follows:

$$\begin{aligned}
& \rho \sum_{s=1}^{N-1} y^s \sum_{k=0}^{s-1} v_{N-1-k} \frac{U_{N-k-2}^{(s-k-1)}(t_0)(-\rho_2)^{s-k-1}}{(s-k-1)!2^{s-k-1}\rho_1^{(N-2-2k+s)/2}} \\
& - \rho_1(1+\alpha\rho) \sum_{s=1}^N y^s \sum_{k=0}^{s-1} v_{N-1-k} \frac{U_{N-k-1}^{(s-k-1)}(t_0)(-\rho_2)^{s-k-1}}{(s-k-1)!2^{s-k-1}\rho_1^{(N-2k+s-1)/2}} \\
& + \rho_1\alpha \sum_{s=0}^{N-1} y^s \sum_{k=0}^s v_{N-1-k} \frac{U_{N-k-1}^{(s-k)}(t_0)(-\rho_2)^{s-k}}{(s-k)!2^{s-k}\rho_1^{(N-2k+s)/2}} + \\
& \left( \rho \sum_{s=0}^{N-1} y^s \frac{U_{N-1}^{(s)}(t_0)(-\rho_2)^s}{(s)!2^s\rho_1^{(N+s)/2}} - \rho_1 \sum_{s=0}^N y^s \frac{U_N^{(s)}(t_0)(-\rho_2)^s}{(s)!2^s\rho_1^{(N+s+1)/2}} \right) p(N, N) + \rho\rho_1^{1/2}p_0 = 0.
\end{aligned}$$

Next we use the relation between the derivatives of the Chebyshev polynomials and Gegenbauer polynomials [3, v.2, p.186]

$$U_n^{(m)}(x) = 2^m m! C_{n-m}^{m+1}(x)$$

to get

$$\begin{aligned}
& \rho \sum_{s=1}^{N-1} y^s \sum_{k=0}^{s-1} v_{N-1-k} \frac{C_{N-s-1}^{s-k}(t_0)(-\rho_2)^{s-k-1}}{\rho_1^{(N-2-2k+s)/2}} \\
& - \rho_1(1+\alpha\rho) \sum_{s=1}^N y^s \sum_{k=0}^{s-1} v_{N-1-k} \frac{C_{N-s}^{s-k}(t_0)(-\rho_2)^{s-k-1}}{\rho_1^{(N-2k+s-1)/2}} \\
& + \rho_1\alpha \sum_{s=0}^{N-1} y^s \sum_{k=0}^s v_{N-1-k} \frac{C_{N-s-1}^{s-k+1}(t_0)(-\rho_2)^{s-k}}{\rho_1^{(N-2k+s)/2}} + \\
& \left( \rho \sum_{s=0}^{N-1} y^s \frac{C_{N-s-1}^{s+1}(t_0)(-\rho_2)^s}{\rho_1^{(N+s)/2}} - \rho_1 \sum_{s=0}^N y^s \frac{C_{N-s}^{s+1}(t_0)(-\rho_2)^s}{\rho_1^{(N+s+1)/2}} \right) p(N, N) + \rho\rho_1^{1/2}p_0 = 0.
\end{aligned}$$

Collecting the terms with the same power of  $y$ , we obtain the required system of equations:

- $s = 0$

$$\alpha\rho_1 C_{N-1}^1(t_0)v_{N-1} + \left[ \rho C_{N-1}^1(t_0) - \rho_1^{1/2} C_N^1(t_0) \right] p(N, N) + \rho\rho_1^{(N+1)/2} p_0 = 0,$$

- $0 < s < N$

$$\begin{aligned} & \rho \sum_{k=0}^{s-1} \frac{C_{N-s-1}^{s-k}(t_0) \rho_1^{k+1}}{(-\rho_2)^{k+1}} v_{N-1-k} - \rho_1^{3/2} (1 + \alpha \rho) \sum_{k=0}^{s-1} \frac{C_{N-s}^{s-k}(t_0) \rho_1^k}{(-\rho_2)^{k+1}} v_{N-1-k} \\ & \quad + \alpha \rho_1 \sum_{k=0}^s \frac{C_{N-s-1}^{s-k+1}(t_0) \rho_1^k}{(-\rho_2)^k} v_{N-1-k} \\ & \quad + \left[ \rho C_{N-s-1}^{s+1}(t_0) - \rho_1^{1/2} C_{N-s}^{s+1}(t_0) \right] p(N, N) = 0, \end{aligned}$$

- $s = N$

$$-\rho_1^{3/2} (1 + \alpha \rho) \sum_{k=0}^{N-1} \frac{C_0^{N-k}(t_0) \rho_1^k}{(-\rho_2)^{k+1}} v_{N-1-k} - \rho_1^{1/2} C_0^{N+1}(t_0) p(N, N) = 0,$$

or, equivalently,

- $s = 0$

$$\alpha \rho_1 C_{N-1}^1(t_0) v_{N+1} + \left[ \rho C_{N-1}^1(t_0) - \rho_1^{1/2} C_N^1(t_0) \right] p(N, N) + \rho \rho_1^{(N+1)/2} p_0 = 0,$$

- $0 < s < N$

$$\begin{aligned} & \sum_{k=0}^{s-1} \left[ \rho \frac{C_{N-s-1}^{s-k}(t_0) \rho_1^{k+1}}{(-\rho_2)^{k+1}} - \rho_1^{3/2} (1 + \alpha \rho) \frac{C_{N-s}^{s-k}(t_0) \rho_1^k}{(-\rho_2)^{k+1}} \right. \\ & \quad \left. + \rho_1 \alpha \frac{C_{N-s-1}^{s-k+1}(t_0) \rho_1^k}{(-\rho_2)^k} \right] v_{N-1-k} + \alpha \frac{C_{N-s-1}^1(t_0) \rho_1^{s+1}}{(-\rho_2)^s} v_{N-1-s} \\ & \quad + \left[ \rho C_{N-s-1}^{s+1}(t_0) - \rho_1^{1/2} C_{N-s}^{s+1}(t_0) \right] p(N, N) = 0, \end{aligned}$$

- $s = N$

$$-\rho_1^{3/2} (1 + \alpha \rho) \sum_{k=0}^{N-1} \frac{C_0^{N-k}(t_0) \rho_1^k}{(-\rho_2)^{k+1}} v_{N-1-k} - \rho_1^{1/2} C_0^{N+1}(t_0) p(N, N) = 0.$$

Finally, to obtain an expression for  $A(y)$  in terms of  $p(k, N)$ ,  $k = 0, \dots, N$  and Chebyshev polynomials, we again substitute subsequently the roots  $a$  and  $b$  into (13) and subtract one equation from another

$$y^N \sum_{k=0}^{N-1} \frac{a^k - b^k}{y^k} p(k, N) - \frac{\rho_1}{\rho_2} \alpha \rho y^{N-1} \sum_{k=0}^{N-1} \frac{a^{k+1} - b^{k+1}}{y^k} p(k, N) \\ + \frac{\rho_1}{\rho_2} \alpha y^{N-1} \sum_{k=0}^{N-1} \frac{a^k - b^k}{y^k} p(k, N) + (a^N - b^N) p(N, N) - \frac{\rho_1}{\rho_2} A(y)(a - b) = 0.$$

As above, taking into account that

$$\frac{a^k - b^k}{a - b} = \frac{U_{k-1}(t)}{\rho_1^{(k-1)/2}},$$

we can express  $A(y)$  in terms of  $p(k, N)$ ,  $k = 0, \dots, N$  and the Chebyshev polynomials of the second type.

## 5. On the numerical calculation of mean queue values

Now we consider the calculation of mean queues. The total mean queue of preemptive and non-preemptive priority packets is given by

$$\bar{n} = \sum_{n=0}^N n \sum_{i=0}^n p(i, n) = \Phi'_y(1, 1) + NV_{N-1}(1) + Np(N, N).$$

The mean queue of preemptive priority packets is equal to

$$\bar{i} = \sum_{n=0}^N \sum_{i=0}^n ip(i, n) = \Phi'_x(1, 1) + V'_{N-1}(1) + Np(N, N).$$

The mean queue of non-preemptive priority packets is  $\bar{n} - \bar{i}$ .

We have also  $V'_{N-1}(1) = \sum_{i=0}^{N-1} ip(i, N)$ .

To derive the unknown functions, we have to do the following steps:

- Substitute  $x = 1$  into the expression for the generating function(12) and reduce it by  $(y - 1)$ , hence

$$(\rho y - 1)\Phi(1, y) = y^N V_{N-1}(1) + y^N p(N, N) - \rho p_0$$

Differentiating by  $y$  and substituting  $y = 1$ , we have:

$$(\rho - 1)\Phi(1, 1) = V_{N-1}(1) + p(N, N) - \rho p_0,$$

$$\rho\Phi(1, 1) + (\rho - 1)\Phi'_y(1, 1) = NV_{N-1}(1) + Np(N, N)$$

- Substitute  $y = 1$  into the expression for the generating function(12) and reduce it by  $(x - 1)$ , hence

$$(\rho_1 x - 1)\Phi(x, 1) = (1 - \alpha\rho_1 x)V_{N-1}(x) + x^N p(N, N) - A(1) - \rho p_0$$

Differentiating by  $x$  and substituting  $x = 1$ , we get:

$$(\rho_1 - 1)\Phi'_x(1, 1) + \rho_1\Phi(1, 1) = -\alpha\rho_1 V_{N-1}(1) + (1 - \alpha\rho_1)V'_{N-1}(1) + Np(N, N)$$

As  $\bar{n} = \Phi'_y(1, 1) + NV_{N-1}(1) + Np(N, N)$ ,  $\bar{i} = \Phi'_x(1, 1) + V'_{N-1}(1) + Np(N, N)$  and the values  $p(i, N)$  are determined by the solution of the system of equations from Theorem 1, the mean queue values are easily determined numerically.

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