

Serdica J. Computing **3** (2009), 159–178

Serdica
Journal of Computing

Bulgarian Academy of Sciences
Institute of Mathematics and Informatics

CAPS IN \mathbb{Z}_n^2

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ABSTRACT. We consider point sets in \mathbb{Z}_n^2 where no three points are on a line – also called caps or arcs. For the determination of caps with maximum cardinality and complete caps with minimum cardinality we provide integer linear programming formulations and identify some values for small n .

1. Introduction. A k -cap in \mathbb{F}_q^d is a subset $A \subseteq \mathbb{F}_q^d$ of size k , where no three points are collinear. A cap is complete if it is not contained in any larger cap. From the combinatorial point of view there are two very natural questions: What is the maximum or minimum size of a cap in \mathbb{F}_q^d ? If we consider caps in projective spaces over \mathbb{F}_q instead of affine spaces then caps are related to linear codes of minimum distance 4, see e. g. [1].

For the maximum size $m_2(\mathbb{F}_q^d)$ of a cap in \mathbb{F}_q^d the following values are known [2, 5, 6]:

$$(1) \quad m_2(\mathbb{F}_q^2) = \begin{cases} q + 2 & \text{for even } q, \\ q + 1 & \text{for odd } q. \end{cases}$$

ACM Computing Classification System (1998): F.2.2.

Key words: Caps, arcs, affine geometry, collinearity, integer programming, rings, complete caps.

- (2) $m_2(\mathbb{F}_q^3) = q^2$ for $q > 2$.
- (3) $m_2(\mathbb{F}_2^d) = 2^d$.
- (4) $m_2(\mathbb{F}_3^4) = 20$, $m_2(\mathbb{F}_3^5) = 45$, and $m_2(\mathbb{F}_4^4) = 40$.

By $n_2(\mathbb{F}_q^d)$ we denote the minimum size of a complete cap in \mathbb{F}_q^d . For the projective case in [4] constructive upper bounds are given. In [3] the authors consider permutations which they interpret as point sets over \mathbb{Z}_n^2 and ask for the minimum number of collinear triples in such configurations. In a similar vein we ask for the maximum cardinality $\sigma(\mathbb{Z}_n^2)$ of a cap in \mathbb{Z}_n^2 where every translation of the two axes contains at most one point.

In this article we consider similar questions in \mathbb{Z}_n^2 instead of \mathbb{F}_q^d . For dimension 2 some authors use the term arcs instead of caps. For p being the smallest prime divisor of an integer n the bounds

$$\max \left\{ 4, \sqrt{2p} + \frac{1}{2} \right\} \leq n_2(\mathbb{Z}_n^2) \leq \max \{4, p + 1\}$$

were proven in [13]. For coprime integers $n, m > 1$ the bound

$$m_2(\mathbb{Z}_{nm}^2) \leq \min \{n \cdot m_2(\mathbb{Z}_m^2), m_2(\mathbb{Z}_n^2) \cdot m\}$$

can be proven. Whenever the value of n is clear from the context we use the abbreviation $\bar{a} := a + \mathbb{Z}n$ for integers a . By $\sigma(\mathbb{Z}_n^2)$ we denote the maximum cardinality of a cap in \mathbb{Z}_n^2 , where each horizontal line $(\bar{1}, \bar{0}) \cdot \mathbb{Z}_n$ and each vertical line $(\bar{0}, \bar{1}) \cdot \mathbb{Z}_n$ contains at most one point. The last conditions model permutations in some sense, see e. g. [3].

1.1. Related work. The *original* “no-three-in-line” problem, introduced by H. Dudeney in 1917, asks if it is possible to select $2n$ points on the n -by- n grid so that no three points are collinear. Currently it is known that for $\varepsilon > 0$ and a sufficiently large n at least $\left(\frac{3}{2} - \varepsilon\right)$ points can be chosen so that no triple is collinear. Guy conjectures that choosing $\frac{\pi}{\sqrt{3}}n \approx 1.814n$ point is asymptotically the best possible. In [21] an analogous question is treated in the three-dimensional space. The question of the minimum size of a complete cap in projective planes over finite fields was originally posed by B. Segre in the late 1950s. In a more general context some authors consider caps (or arcs) over so-called projective Hjelmslev planes, see e. g. [8, 10, 11, 12, 14, 15, 18, 19]. Here we note that \mathbb{Z}_p^2 is the affine part for the chain ring \mathbb{Z}_{p^r} .

1.2. Our contribution. In this article we develop an algorithm which can decide whether three given points in \mathbb{Z}_n^2 are collinear or not in $O(n \log n)$, given the prime factorization of n . We model the problem of the exact determination of $m_2(\mathbb{Z}_n^2)$, $n_2(\mathbb{Z}_n^2)$, and $\sigma(\mathbb{Z}_n^2)$ as integer linear programs. Finally we perform some computer calculations to determine some as yet unknown values, e. g. we validate $m_2(\mathbb{Z}_{25}^2) = 20$.

1.3. Organization of the paper. In Section 2 we specify when we consider three points of \mathbb{Z}_n^2 to be collinear and develop a fast algorithm which can check collinearity. To this end some cumbersome and technical but elementary calculations have to be executed. In Section 3 we give integer linear programming formulations and in Section 4 we combine them with some symmetry-breaking techniques to determine the exact values of $m_2(\mathbb{Z}_n^2)$, $n_2(\mathbb{Z}_n^2)$, and $\sigma(\mathbb{Z}_n^2)$ for small n .

2. Points on a line. A line in \mathbb{Z}_n^2 is a translate of a cyclic subgroup of order n . We remark that every cyclic subgroup of \mathbb{Z}_n^2 is contained in some subgroup of order n , see also [13]. An example is given by the line $(\bar{3}, \bar{7}) + (\bar{1}, \bar{5}) \cdot \mathbb{Z}_{12}$ in \mathbb{Z}_{12}^2 , see Figure 1. A point p is called **incident** with a line l if $p \in l$. With this we could define r points to be **collinear** if they are incident with a common line.

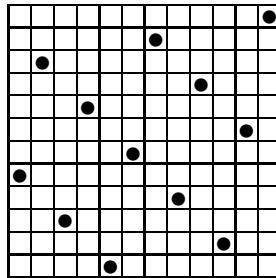


Fig. 1. The line $(\bar{3}, \bar{7}) + (\bar{1}, \bar{5}) \cdot \mathbb{Z}_{12}$ over \mathbb{Z}_{12}^2 .

Lemma 2.1. *A set of r points $(u_i, v_i) \in \mathbb{Z}_n^2$ is collinear if there exist $a, b, t_1, t_2, w_i \in \mathbb{Z}_n$ with*

$$a + w_i t_1 = u_i \quad \text{and} \quad b + w_i t_2 = v_i$$

for $1 \leq i \leq r$.

Proof. We can write a line L in \mathbb{Z}_n^2 as a set

$$\left\{ (a + wt_1, b + wt_2) \mid w \in \mathbb{Z}_n \right\},$$

where a, b, t_1, t_2 are arbitrary elements of \mathbb{Z}_n . \square

We would like to remark that one could also define a line as the set of solutions $(x, y) \in \mathbb{Z}_n^2$ of where $ax + by = c$, where $a, b, c \in \mathbb{Z}_n$ with $\gcd(\widehat{a}, \widehat{b}, n) = 1$, see [13].

If $n = p$ is a prime then every two points of \mathbb{Z}_p^2 uniquely determine a line. This does not hold in general for arbitrary n . If $n = p^r$ is a prime power then every two points uniquely determine p^k lines containing those points, where k depends on the points, see e. g. [10, 19]. It is possible to define neighborhood relations \sim_k for points a, b by requiring that there exist at least p^k lines containing a and b . For arbitrary n two distinct points are incident with at least one line.

If $n = p$ is a prime then \mathbb{Z}_n is a field and there is a well-known test to check whether three points are collinear or not, which runs in time $O(1)$:

Lemma 2.2. *For a prime n three points $(u_1, v_1), (u_2, v_2), (u_3, v_3) \in \mathbb{Z}_n^2$ are collinear if and only if*

$$(1) \quad \begin{vmatrix} u_1 & v_1 & \bar{1} \\ u_2 & v_2 & \bar{1} \\ u_3 & v_3 & \bar{1} \end{vmatrix} = \bar{0}.$$

We note that in \mathbb{Z}_8^2 the points $(\bar{0}, \bar{0}), (\bar{2}, \bar{4}), (\bar{4}, \bar{4})$ fulfill Equation (1) but are not collinear. So in general equation (1) is necessary but not sufficient for three points to be collinear.

Lemma 2.3. *If three points $(u_1, v_1), (u_2, v_2), (u_3, v_3) \in \mathbb{Z}_n^2$ are collinear, then*

$$\begin{vmatrix} u_1 & v_1 & \bar{1} \\ u_2 & v_2 & \bar{1} \\ u_3 & v_3 & \bar{1} \end{vmatrix} = \bar{0}$$

holds.

Proof. Due to Lemma 2.1 we have

$$\begin{vmatrix} u_1 & v_1 & \bar{1} \\ u_2 & v_2 & \bar{1} \\ u_3 & v_3 & \bar{1} \end{vmatrix} = \begin{vmatrix} a + w_1t_1 & b + w_1t_2 & \bar{1} \\ a + w_2t_1 & b + w_2t_2 & \bar{1} \\ a + w_3t_1 & b + w_3t_2 & \bar{1} \end{vmatrix} = \begin{vmatrix} w_1t_1 & w_1t_2 & \bar{1} \\ w_2t_1 & w_2t_2 & \bar{1} \\ w_3t_1 & w_3t_2 & \bar{1} \end{vmatrix} = \bar{0}. \quad \square$$

In the following we will develop an algorithm which can decide if three given points in \mathbb{Z}_n^2 are collinear in time $O\left(\sum_{i=1}^l \mu_i\right)$ using the prime factorization of $n = \prod_{i=1}^l p_i^{\mu_i}$ as input. We would like to remark that also the Smith normal form can be utilized to obtain an $O(\log n)$ -algorithm.

At first we remark that w.l.o.g. we can assume that one of three points equals $(\bar{0}, \bar{0})$:

Lemma 2.4. *Three points $(u_1, v_1), (u_2, v_2), (u_3, v_3) \in \mathbb{Z}_n^2$ are collinear, if and only if there exist $t_1, t_2, w_2, w_3 \in \mathbb{Z}_n$ with*

$$\begin{aligned} w_2 t_1 &= u_2 - u_1, \\ w_3 t_1 &= u_3 - u_1, \\ w_2 t_2 &= v_2 - v_1, \text{ and} \\ w_3 t_2 &= v_3 - v_1. \end{aligned}$$

Proof. Since $(u_1, v_1), (u_2, v_2), (u_3, v_3)$ are collinear if and only if $(0, 0), (u_2 - u_1, v_2 - v_1), (u_3 - u_1, v_3 - v_1)$ are collinear, there exist $a', b', t'_1, t'_2, w'_1, w'_2, w'_3 \in \mathbb{Z}_n$ with

$$\begin{aligned} a' + w'_1 t'_1 &= \bar{0}, \\ a' + w'_2 t'_1 &= u_2 - u_1, \\ a' + w'_3 t'_1 &= u_3 - u_1, \\ b' + w'_1 t'_2 &= \bar{0}, \\ b' + w'_2 t'_2 &= v_2 - v_1, \text{ and} \\ b' + w'_3 t'_2 &= v_3 - v_1. \end{aligned}$$

If we have a solution $t_1, t_2, w_2, w_3 \in \mathbb{Z}_n$ of the first equation system then $a' = b' = w'_1 = 0, t'_1 = t_1, t'_2 = t_2, w'_2 = w_2, w'_3 = w_3$ is a solution of the second equation system.

If we have a solution $a', b', t'_1, t'_2, w'_1, w'_2, w'_3 \in \mathbb{Z}_n$ of the second equation system then $a' = -w'_1 t'_1$ and $b' = -w'_1 t'_2$ holds. With this $t_1 = t'_1, t_2 = t'_2, w_2 = w'_2 - w'_1, w_3 = w'_3 - w'_1$ is a solution of the first equation system. \square

Due to the Chinese remainder theorem we have:

Lemma 2.5. *If $n = a \cdot b$ with coprime a and b , then three points $p_1, p_2, p_3 \in \mathbb{Z}_n^2$ are collinear if and only if both projections into \mathbb{Z}_a^2 and \mathbb{Z}_b^2 give a collinear point set.*

Thus it suffices to consider the case where $n = p^r$ is a prime power. For $a \in \mathbb{Z}_n$ let $\hat{a} \in \{1, \dots, n\}$ denote the integer fulfilling $\hat{a} + \mathbb{Z}_n = a$. If none of the values u_2, u_3, v_2, v_3 is invertible in \mathbb{Z}_{p^r} we can apply the following reduction:

Lemma 2.6. *If for $r \geq 2$ the prime p divides $\hat{u}_2, \hat{u}_3, \hat{v}_2$, and \hat{v}_3 then the points $(\bar{0}, \bar{0}), (u_2, v_2)$ and (u_3, v_3) are collinear in $\mathbb{Z}_{p^r}^2$ if and only if the points*

$$(\bar{0}, \bar{0}), \left(\frac{\hat{u}_2}{p} + \mathbb{Z}p^{r-1}, \frac{\hat{v}_2}{p} + \mathbb{Z}p^{r-1} \right), \text{ and } \left(\frac{\hat{u}_3}{p} + \mathbb{Z}p^{r-1}, \frac{\hat{v}_3}{p} + \mathbb{Z}p^{r-1} \right)$$

are collinear in $\mathbb{Z}_{p^{r-1}}^2$.

Proof. We have the following equivalence: If and only if the points $(\bar{0}, \bar{0}), (u_2, v_2)$ and (u_3, v_3) are collinear in $\mathbb{Z}_{p^r}^2$ then due to Lemma 2.4 there exist integers $\tilde{t}_1, \tilde{t}_2, \tilde{w}_2, \tilde{w}_3 \in \{1, \dots, p^r\}$ fulfilling

$$p^r \mid \hat{u}_2 - \tilde{w}_2 \tilde{t}_1, \quad p^r \mid \hat{u}_3 - \tilde{w}_3 \tilde{t}_1, \quad p^r \mid \hat{v}_2 - \tilde{w}_2 \tilde{t}_2, \quad \text{and } p^r \mid \hat{v}_3 - \tilde{w}_3 \tilde{t}_2.$$

Similar we conclude: If and only if the points

$$(\bar{0}, \bar{0}), \left(\frac{\hat{u}_2}{p} + \mathbb{Z}p^{r-1}, \frac{\hat{v}_2}{p} + \mathbb{Z}p^{r-1} \right), \text{ and } \left(\frac{\hat{u}_3}{p} + \mathbb{Z}p^{r-1}, \frac{\hat{v}_3}{p} + \mathbb{Z}p^{r-1} \right)$$

are collinear in $\mathbb{Z}_{p^{r-1}}^2$ then there exist integers $\dot{t}_1, \dot{t}_2, \dot{w}_2, \dot{w}_3 \in \{1, \dots, p^{r-1}\}$ fulfilling

$$p^{r-1} \mid \frac{\hat{u}_2}{p} - \dot{w}_2 \dot{t}_1, \quad p^{r-1} \mid \frac{\hat{u}_3}{p} - \dot{w}_3 \dot{t}_1, \quad p^{r-1} \mid \frac{\hat{v}_2}{p} - \dot{w}_2 \dot{t}_2, \quad \text{and } p^{r-1} \mid \frac{\hat{v}_3}{p} - \dot{w}_3 \dot{t}_2.$$

If the tuple $\dot{t}_1, \dot{t}_2, \dot{w}_2, \dot{w}_3 \in \{1, \dots, p^{r-1}\}$ is a solution of the second system, then $\tilde{t}_1 = \dot{t}_1 \cdot p, \tilde{t}_2 = \dot{t}_2 \cdot p, \tilde{w}_2 = \dot{w}_2, \tilde{w}_3 = \dot{w}_3$ is a solution with $\tilde{t}_1, \tilde{t}_2, \tilde{w}_2, \tilde{w}_3 \in \{1, \dots, p^r\}$ of the first system.

If the tuple $\tilde{t}_1, \tilde{t}_2, \tilde{w}_2, \tilde{w}_3 \in \{1, \dots, p^r\}$ is a solution of the first system then $\dot{t}_1 = \frac{\tilde{t}_1}{p}, \dot{t}_2 = \frac{\tilde{t}_2}{p}, \dot{w}_2 = \tilde{w}_2, \dot{w}_3 = \tilde{w}_3$ is a solution with $\dot{t}_1, \dot{t}_2, \dot{w}_2, \dot{w}_3 \in \{1, \dots, p^{r-1}\}$ of the second system. \square

Thus in the following we can confine ourselves to the case where $r \geq 2$ (for $r = 1$ we can utilize Lemma 2.2) and at least one of u_2, u_3, v_2 , or v_3 is invertible in \mathbb{Z}_{p^r} . We claim that in such a situation the criterion of Lemma 2.3 is also sufficient for $(\bar{0}, \bar{0}), (u_2, v_2)$, and (u_3, v_3) being collinear:

Lemma 2.7. *If at least one of the elements $u_2, u_3, v_2, v_3 \in \mathbb{Z}_{p^r}$ is invertible and Equation (1) is fulfilled then the three points $(\bar{0}, \bar{0}), (u_2, v_2), (u_3, v_3)$ in $\mathbb{Z}_{p^r}^2$ are collinear.*

Proof. Due to symmetry we can assume that u_2 is invertible. Setting $t_1 = \bar{1}$, $w_2 = u_2$, $t_2 = v_2 u_2^{-1}$, and $w_3 = u_3$, using the notation of Lemma 2.4, we obtain $v_3 = u_3 v_2 u_2^{-1}$. Since this equation is equivalent to Equation (1) we can use Lemma 2.4 to conclude that $(\bar{0}, \bar{0})$, (u_2, v_2) , and (u_3, v_3) are collinear. \square

Using the previous lemmas we can design an efficient algorithm, with some subroutines, to decide whether $r \geq 3$ points $(u_i, v_i) \in \mathbb{Z}_n^2$ are collinear or not. We assume that the prime factorization

$$n = \prod_{i=1}^l p_i^{\mu_i}$$

is known in advance. For practical purposes our algorithm deals with integers instead of residue classes.

Algorithm 2.8. `is_collinear`($\hat{u}_1, \hat{v}_1, \dots, \hat{u}_r, \hat{v}_r, \mathbf{p}_1, \mu_1, \dots, \mathbf{p}_l, \mu_l$)

```

n =  $\prod_{i=1}^l p_i^{\mu_i}$ 
 $u_2 = \hat{u}_2 - \hat{u}_1, v_2 = \hat{v}_2 - \hat{v}_1$ 
if  $u_2 \leq 0$  then  $u_2 = u_2 + n$ 
if  $v_2 \leq 0$  then  $v_2 = u_2 + n$ 
for  $i$  from 3 to  $r$ 
   $u_3 = \hat{u}_i - \hat{u}_1, v_3 = \hat{v}_i - \hat{v}_1$ 
  if  $u_3 \leq 0$  then  $u_3 = u_3 + n$ 
  if  $v_3 \leq 0$  then  $v_3 = u_3 + n$ 
  if is_collinear_fix_zero( $u_2, v_2, u_3, v_3, \mathbf{p}_1, \mu_1, \dots, \mathbf{p}_l, \mu_l$ ) == false
    then return false
return true

```

Algorithm 2.9. `is_collinear_fix_zero`($u_2, v_2, u_3, v_3, \mathbf{p}_1, \mu_1, \dots, \mathbf{p}_l, \mu_l$)

```

for  $i$  from 1 to  $l$ 
   $k = p_i^{\mu_i}$ 
   $u'_2 = u_2 - \left\lfloor \frac{u_2}{k} \right\rfloor \cdot k$ 
   $u'_3 = u_3 - \left\lfloor \frac{u_3}{k} \right\rfloor \cdot k$ 
   $v'_2 = v_2 - \left\lfloor \frac{v_2}{k} \right\rfloor \cdot k$ 
   $v'_3 = v_3 - \left\lfloor \frac{v_3}{k} \right\rfloor \cdot k$ 
  if is_collinear_prime_power( $u'_2, v'_2, u'_3, v'_3, \mathbf{p}_i, \mu_i$ ) == false

```

then return *false*
return *true*

Algorithm 2.10. `is_collinear_prime_power`($\mathbf{u}_2, \mathbf{v}_2, \mathbf{u}_3, \mathbf{v}_3, \mathbf{p}, r$)
if $r == 1$ then
 if $u_2 v_3 \equiv u_3 v_2 \pmod{p}$ then return *true* else return *false*
else
 if $u_2 \equiv u_3 \equiv v_2 \equiv v_3 \equiv \bar{0} \pmod{p}$
 then return `is_collinear_prime_power` $\left(\frac{\mathbf{u}_2}{\mathbf{p}}, \frac{\mathbf{v}_2}{\mathbf{p}}, \frac{\mathbf{u}_3}{\mathbf{p}}, \frac{\mathbf{v}_3}{\mathbf{p}}, \mathbf{p}, r - 1\right)$
 else
 if $u_2 v_3 \equiv u_3 v_2 \pmod{p^r}$ then return *true* else return *false*

Now we want to analyze the running time of Algorithm 2.8 and Algorithm 2.9:

Theorem 2.11. *If $n = \prod_{i=1}^l p_i^{\mu_i}$ then Algorithm 2.9 needs at most $O\left(\sum_{i=1}^l \mu_i\right)$ time steps.*

Proof. It suffices to prove that Algorithm 2.10 needs at most r recursions, which is obvious. \square

Corollary 2.12. *Given the prime factorization of n Algorithm 2.9 runs in $O(\log n)$ and Algorithm 2.8 runs in $O(r \cdot \log n)$ time.*

So let us take a (small) example to illustrate Algorithm 2.8. We choose $n = 625 = 5^4$, $u_1 = \bar{1}$, $v_1 = \bar{2}$, $u_2 = \bar{76}$, $v_2 = \bar{57}$, $u_3 = \bar{251}$ and $v_3 = \bar{102}$. At first we transform the problem to $u_1 = v_1 = \bar{0}$, $u_2 = \bar{75}$, $v_2 = \bar{55}$, $u_3 = \bar{250}$, and $v_3 = \bar{100}$. Since n is a prime power we do not split up the problem into prime powers. Since the largest power of 5 which divides all of \tilde{u}_2 , \tilde{u}_3 , \tilde{v}_2 , and \tilde{v}_3 is 5^1 we reduce the problem to $(\bar{0}, \bar{0})$, $(\bar{15}, \bar{11})$, $(\bar{50}, \bar{20})$ in \mathbb{Z}_{125} . Due to $\bar{15} \cdot \bar{20} = \bar{11} \cdot \bar{50}$ in \mathbb{Z}_{125} the three original points are collinear.

If we have to check very often whether three points are collinear or not then it is more efficient to create a $\mathbb{Z}_n^2 \times \mathbb{Z}_n^2$ table in order to bookmark whether $(\bar{0}, \bar{0})$, p_1 , p_2 are collinear or not.

3. Integer Linear Programming formulations. In this section we formulate the problem of the exact determination of $m_2(\mathbb{Z}_n^2)$ as an integer linear program using the binary variables $x_{i,j} \in \{0, 1\}$ for $1 \leq i, j \leq n$. Here the variables $x_{i,j}$ encode a subset

$$C := \left\{ (i + \mathbb{Z}_n, j + \mathbb{Z}_n) \mid x_{i,j} = 1, 1 \leq i, j \leq n \right\} \subseteq \mathbb{Z}_n^2.$$

To enforce C to be a cap, we require the linear inequality

$$(2) \quad \sum_{i,j : (i+\mathbb{Z}_n, j+\mathbb{Z}_n) \in L} x_{i,j} \leq 2$$

for all lines L of \mathbb{Z}_n^2 . It is not difficult to show that these inequalities suffice to enforce that no three points of C are collinear. Obviously we could also write up an inequality for every triple of collinear points, but we note that Inequality (2) is more compact. With the above we can state

$$(3) \quad m_2(\mathbb{Z}_n^2) = \max \sum_{i=1}^n \sum_{j=1}^n x_{i,j}$$

subject to

$$\sum_{i,j : (i+\mathbb{Z}_n, j+\mathbb{Z}_n) \in L} x_{i,j} \leq 2 \quad \forall \text{ lines } L \text{ of } \mathbb{Z}_n^2$$

$$x_{i,j} \in \{0, 1\} \quad \forall 1 \leq i, j \leq n.$$

We note that every optimal solution of ILP (3) corresponds to a cap which is complete.

Similarly we can state for the maximum cardinality of a cap which is a subset of a permutation point set:

$$(4) \quad \sigma(\mathbb{Z}_n^2) = \max \sum_{i=1}^n \sum_{j=1}^n x_{i,j}$$

subject to

$$\sum_{i,j : (i+\mathbb{Z}_n, j+\mathbb{Z}_n) \in L} x_{i,j} \leq 2 \quad \forall \text{ lines } L \text{ of } \mathbb{Z}_n^2$$

$$\sum_{i=1}^n x_{i,j} \leq 1 \quad \forall 1 \leq j \leq n$$

$$\sum_{j=1}^n x_{i,j} \leq 1 \quad \forall 1 \leq i \leq n$$

$$x_{i,j} \in \{0, 1\} \quad \forall 1 \leq i, j \leq n.$$

Some more work is needed to express $n_2(\mathbb{Z}_n^2)$ as the optimal objective value of an ILP. Simply replacing max by min in ILP (3) would yield the optimal solution

$x_{i,j} = 0$ for all $1 \leq i, j \leq n$. So we have to augment ILP (3) by some additional conditions and variables in order to enforce the set C , corresponding to $x_{i,j}$ to be complete. Therefore we introduce line variables $y_L \in \{0, 1\}$ for every line L in \mathbb{Z}_n^2 . The idea is that y_L should equal 1 if C contains exactly two points of L . This can be modeled using the linear inequality

$$(5) \quad 1 + y_L \geq \sum_{i,j:(i+\mathbb{Z}_n,j+\mathbb{Z}_n) \in L} x_{i,j} \geq 2y_L$$

for all lines L in \mathbb{Z}_n^2 .

To model the completeness of C we introduce the linear inequality

$$(6) \quad x_{i,j} + \sum_{L:(i+\mathbb{Z}_n,j+\mathbb{Z}_n) \in L} y_L \geq 1$$

for all $1 \leq i, j \leq n$. Here the idea is, that a cap C is complete if and only if there does not exist a point $P \in \mathbb{Z}_n^2 \setminus C$ such that $C \cup \{P\}$ is also a cap. So let us assume that we have a binary variable allocation $x_{i,j}, y_L$ satisfying Inequalities (2), (5), and (6), then for all $1 \leq i, j \leq n$ either $P = (i + \mathbb{Z}_n, j + \mathbb{Z}_n)$ is contained in C or there exist two points P_1, P_2 in C such that P_1, P_2 , and P are collinear. Thus every feasible solution of the ILP

$$(7) \quad n_2(\mathbb{Z}_n^2) = \min \sum_{i=1}^n \sum_{j=1}^n x_{i,j}$$

subject to

$$\begin{aligned} \sum_{i,j:(i+\mathbb{Z}_n,j+\mathbb{Z}_n) \in L} x_{i,j} &\leq 2 && \forall \text{ lines } L \text{ of } \mathbb{Z}_n^2 \\ \sum_{i,j:(i+\mathbb{Z}_n,j+\mathbb{Z}_n) \in L} x_{i,j} - 2y_L &\geq 0 && \forall \text{ lines } L \text{ of } \mathbb{Z}_n^2 \\ x_{i,j} + \sum_{L:(i+\mathbb{Z}_n,j+\mathbb{Z}_n) \in L} y_L &\geq 1 && \forall 1 \leq i, j \leq n \\ x_{i,j} &\in \{0, 1\} && \forall 1 \leq i, j \leq n \\ y_L &\in \{0, 1\} && \forall \text{ lines } L \text{ of } \mathbb{Z}_n^2 \end{aligned}$$

corresponds to a complete cap C . For a given complete cap C we can extend the corresponding partial variable allocation by setting $y_L = 1$ exactly if C contains exactly two points of C . Since we minimize the number of points of C the target value of ILP (7) equals $n_2(\mathbb{Z}_n^2)$.

Now let us have a look at the number of variables and inequalities of the ILPs (3), (4), and (7). A function $\psi : \mathbb{N} \rightarrow \mathbb{N}$ is called **multiplicative** if $\psi(1) = 1$ and $\psi(nm) = \psi(n) \cdot \psi(m)$ for all coprime integers n and m .

Definition 3.1. Let $\psi : \mathbb{N} \rightarrow \mathbb{N}$ be the multiplicative arithmetic theoretic function defined by $\psi(p^r) = (p+1)p^{r-1}$ for prime powers $p^r > 1$.

We note that since

$$\lim_{k \rightarrow \infty} \ln p_k \prod_{i=1}^k \frac{1}{1 + \frac{1}{p_i}} = \frac{\pi^2}{6e^\gamma},$$

and

$$\lim_{k \rightarrow \infty} \left(\prod_{i=1}^k p_i \right)^{\frac{1}{p_k}} = e$$

where p_k denotes the k th prime and γ denotes the Euler-Mascheroni constant with an approximate value of 0.57721566, and due to $\psi(n^2) = n^2 \cdot \prod_{p|n} 1 + \frac{1}{p}$ (for primes p), we have

$$n^2 \leq \psi(n^2) \leq 1.0828 \cdot n^2 \ln \ln n$$

for all sufficiently large n .

Lemma 3.2. There are $\psi(n^2)$ lines in \mathbb{Z}_n^2 .

Proof. Due to the Chinese remainder theorem the number of lines is a multiplicative arithmetic function. The number of cyclic subgroups of order p^r in $\mathbb{Z}_{p^r}^2$ is given by $2\varphi(p^r)(p^r - \varphi(p^r)) + \varphi(p^r)^2 = p^{2r} - p^{2r-2}$, where $\varphi(p^r) = p^{r-1}(p-1)$ is Euler's totient function. Since every cyclic subgroups of order p^r contains $\varphi(p^r) = p^{r-1}(p-1)$ generators, there are $\frac{p^{2r} - p^{2r-2}}{p^{r-1}(p-1)} = (p+1)p^{r-1}$ lines through each point in $\mathbb{Z}_{p^r}^2$. In total there are $\frac{(p+1)p^{r-1} \cdot p^{2r}}{p^r} = \psi(p^{2r})$ lines since every line contains p^r points. \square

Thus ILP (3) consists of n^2 variables and $\psi(n^2)$ inequalities, ILP (4) consists of n^2 variables and $\psi(n^2) + 2n$ inequalities, and ILP (7) consists of $n^2 + \psi(n^2)$ variables and $2\psi(n^2) + n^2$ inequalities. So in all cases the number of variables and inequalities are in $O(n^2 \ln n)$. But since generally the optimization (or also the feasibility problem) of 0-1 linear programs is NP-complete these ILP formulations might not help too much from the theoretical point of view. On the

other hand these ILP formulations enable us to determine some exact numbers and bounds of $m_2(\mathbb{Z}_n^2)$, $n_2(\mathbb{Z}_n^2)$, and $\sigma(\mathbb{Z}_n^2)$ in the next section.

In contrast to ILP problems LP problems, i. e. ILP problems without integrality constraints, can be solved in polynomial time. So in order to obtain an LP for the ILPs (3), (4), and (7) we can relax the conditions $x_{i,j} \in \{0,1\}$, $y_L \in \{0,1\}$ by $0 \leq x_{i,j} \leq 1$, $0 \leq y_L \leq 1$. To solve the original (integral) problem several techniques, e. g. branch & bound, have to be applied. In many cases additional inequalities will be very useful for an optimization algorithm. We will explain this idea by considering the example

$$\begin{aligned} & \max x_1 + 2x_2 \\ & \text{subject to} \\ & 5x_1 + 3x_2 \leq 15 \\ & x_2 \leq 2 \\ & x_1, x_2 \in \mathbb{N}_0 \end{aligned}$$

On the left-hand side of Figure we have depicted the feasible set of the relaxed linear program (i. e. we have replaced $x_1, x_2 \in \mathbb{N}_0$ by $x_1, x_2 \geq 0$). The integral points are marked by filled circles. If we additionally require $x_1 + x_2 \leq 3$ we obtain the feasible set as depicted on the right-hand side of Figure 2.



Fig. 2. Feasible sets of linear programs

We observe that the feasible set on the right-hand side contains the same integral points as the feasible set of the left-hand side, whereas the surface area differs. In this case we say that $x_1 + x_2 \leq 3$ is a valid inequality w. r. t. the integral points. If the right-hand side of Inequality (2) in ILP (3) would be 1 instead of 2, then several classes of valid inequalities of the so-called stable set polytope are known, e. g.. odd circuit inequalities or clique inequalities (if given by edge inequalities).

Unfortunately we are not aware of any general (masking the geometric properties of \mathbb{Z}_n^2) results on valid inequalities for the polytope of ILP (3).

4. Bounds and exact values. In this section we want do state bounds and exact values for the three problems stated in the introduction which are either known or derived using an ILP based approach following the ILP formulations of Section 3.

4.1. Maximum cardinality of caps over \mathbb{Z}_n^2 . As mentioned in the introduction we have:

Lemma 4.1. $m_2(\mathbb{Z}_{nm}^2) \leq \min \{n \cdot m_2(\mathbb{Z}_m^2), m_2(\mathbb{Z}_n^2) \cdot m\}$ for coprime integers $n, m > 1$.

Proof. Let $X \subseteq \mathbb{Z}_{nm}^2$ be a cap, l be a line in \mathbb{Z}_n^2 , and $q_1 : \mathbb{Z}_{nm}^2 \rightarrow \mathbb{Z}_n^2$, $q_2 : \mathbb{Z}_{nm}^2 \rightarrow \mathbb{Z}_m^2$ be reduction maps. We set $Y := \{x \in X \mid q_1(x) \in l\}$. If $|Y| \geq 2$ then $q_1(Y) \subseteq q_1(X)$ is collinear in \mathbb{Z}_n^2 . Thus we must have $|Y| \leq m_2(\mathbb{Z}_m^2)$ since otherwise $q_2(Y)$ is collinear in \mathbb{Z}_m^2 , from which we could conclude the collinearity of Y in \mathbb{Z}_{nm}^2 . Since we can partition \mathbb{Z}_n^2 into n lines and $m_2(\mathbb{Z}_m^2) \geq 2$ we have $|X| \leq n \cdot m_2(\mathbb{Z}_m^2)$. Due to symmetry we also have $|X| \leq m \cdot m_2(\mathbb{Z}_n^2)$. \square

Directly solving ILP (3) using the commercial solver **ILOG CPLEX 11.2** yields the results given in Table 1. We would like to remark that for $n \leq 17$ the exact values of $m_2(\mathbb{Z}_n^2)$ were also determined in [14].

n	2	3	4	5	6	7	8	9	10	11	12
$m_2(\mathbb{Z}_n^2)$	4	4	6	6	8	8	8	9	12	12	12
time in s	< 0.1	< 0.1	< 0.1	< 0.1	< 0.1	0.65	0.12	48.05	5.30	126.24	56.80

Table 1. Directly solving ILP (3).

One of the main reasons why ILP solvers fail to tackle instances of ILP (3) for *larger* values of n is the highly symmetric formulation of the problem, i. e. instead of few optimal solutions there exist a whole bulk of solutions which correspond to geometrically isomorphic caps. So we have to break symmetries by introducing further inequalities. To this end we consider automorphisms α of \mathbb{Z}_n^2 , i. e. mappings from \mathbb{Z}_n^2 to \mathbb{Z}_n^2 which preserve point-line incidences. This means, in particular, that $C \subseteq \mathbb{Z}_n^2$ is a cap if and only if $\alpha(C)$ is a cap. Since the translations $x \mapsto x + y$ for $y \in \mathbb{Z}_n^2$ are such automorphisms we can assume w.l.o.g. that the point $(\bar{0}, \bar{0})$ is contained in each non-empty cap C . Another class of automorphisms arises from multiplication of invertible 2×2 matrices of \mathbb{Z}_n . So let us denote the resulting group of all translations and all invertible 2×2 matrices by G . (All elements of G are automorphisms.)

In the following we will assume that some points $a_i \in \mathbb{Z}_n^2$ are not contained in cap C but some points $b_i \in \mathbb{Z}_n^2$ are contained in C . From this we can conclude

n	cases
14	$a_1 = (\overline{0}, \overline{0}), b_1 = (\overline{1}, \overline{0})$ resolved in 0.57 s. $a_1 = (\overline{0}, \overline{0}), a_2 = (\overline{1}, \overline{0}), b_1 = (\overline{0}, \overline{1})$ resolved in 3.21 s. $a_1 = (\overline{0}, \overline{0}), a_2 = (\overline{1}, \overline{0}), a_3 = (\overline{0}, \overline{1})$ resolved in 675.54 s.
15	$a_1 = (\overline{0}, \overline{0}), b_1 = (\overline{1}, \overline{0})$ resolved in 0.04 s. $a_1 = (\overline{0}, \overline{0}), a_2 = (\overline{1}, \overline{0}), b_1 = (\overline{0}, \overline{1})$ resolved in 0.67 s. $a_1 = (\overline{0}, \overline{0}), a_2 = (\overline{1}, \overline{0}), a_3 = (\overline{0}, \overline{1})$ resolved in 1830.68 s.
16	$a_1 = (\overline{0}, \overline{0}), b_1 = (\overline{1}, \overline{0})$ resolved in 0.17 s. $a_1 = (\overline{0}, \overline{0}), a_2 = (\overline{1}, \overline{0}), b_1 = (\overline{0}, \overline{1})$ resolved in 0.54 s. $a_1 = (\overline{0}, \overline{0}), a_2 = (\overline{1}, \overline{0}), a_3 = (\overline{0}, \overline{1})$ resolved in 251.93 s.
18	$a_1 = (\overline{0}, \overline{0}), b_1 = (\overline{1}, \overline{0})$ resolved in 0.07 s. $a_1 = (\overline{0}, \overline{0}), a_2 = (\overline{1}, \overline{0}), b_1 = (\overline{0}, \overline{1})$ resolved in 0.13 s. $a_1 = (\overline{0}, \overline{0}), a_2 = (\overline{1}, \overline{0}), a_3 = (\overline{0}, \overline{1})$ resolved in 2554.91 s.
20	$a_1 = (\overline{0}, \overline{0}), b_1 = (\overline{1}, \overline{0})$ resolved in less than 1 s. $a_1 = (\overline{0}, \overline{0}), a_2 = (\overline{1}, \overline{0}), b_1 = (\overline{0}, \overline{1})$ resolved in less than 1 s. $a_1 = (\overline{0}, \overline{0}), a_2 = (\overline{1}, \overline{0}), a_3 = (\overline{0}, \overline{1}), a_4 = (\overline{4}, \overline{1})$ resolved in 1950 s. $a_1 = (\overline{0}, \overline{0}), a_2 = (\overline{1}, \overline{0}), a_3 = (\overline{0}, \overline{1}), a_4 = (\overline{4}, \overline{4}), b_1 = (\overline{4}, \overline{1})$ resolved in 513 s. $a_1 = (\overline{0}, \overline{0}), a_2 = (\overline{1}, \overline{0}), a_3 = (\overline{0}, \overline{1}), a_4 = (\overline{5}, \overline{4}), b_1 = (\overline{4}, \overline{1}),$ $b_2 = (\overline{4}, \overline{4})$ resolved in 80 s. $\dots, a_4 = (\overline{5}, \overline{3}), \dots$ resolved in less than 1 s. $\dots, a_4 = (\overline{9}, \overline{1}), \dots$ resolved in less than 1 s. $\dots, a_4 = (\overline{16}, \overline{9}), \dots$ resolved in 3584 s. $\dots, a_4 = (\overline{15}, \overline{11}), \dots$ resolved in 800 s. $\dots, a_4 = (\overline{15}, \overline{10}), \dots$ resolved in 3562 s. $\dots, a_4 = (\overline{6}, \overline{5}), \dots$ resolved in 3725 s. $\dots, a_4 = (\overline{5}, \overline{5}), \dots$ resolved in 5922 s. $\dots, a_4 = (\overline{10}, \overline{1}), \dots$ resolved in less than 1 s. $\dots, a_4 = (\overline{5}, \overline{1}), \dots$ resolved in 3080 s. $a_1 = (\overline{0}, \overline{0}), a_2 = (\overline{1}, \overline{0}), a_3 = (\overline{0}, \overline{1}), b_1 = (\overline{4}, \overline{1}), b_2 = (\overline{4}, \overline{4}), b_3 = (\overline{5}, \overline{4}),$ $b_4 = (\overline{5}, \overline{3}), b_5 = (\overline{9}, \overline{1}), b_6 = (\overline{16}, \overline{9}), b_7 = (\overline{15}, \overline{11}), b_8 = (\overline{15}, \overline{10}), b_9 = (\overline{6}, \overline{5}),$ $b_{10} = (\overline{5}, \overline{5}), b_{11} = (\overline{10}, \overline{1}), b_{12} = (\overline{5}, \overline{1})$ resolved in less than 1 s.

Table 2. Case differentiations for the determination of $m_2(\mathbb{Z}_n^2)$

some valid inequalities for our ILPs using the group G . If there exists an index h , an element $z \in \mathbb{Z}_n^2$, and an automorphism α such that

$$\alpha(\{z, a_1, \dots\}) = \{b_h, a_1, \dots\}$$

then we can conclude w.l.o.g. that $z \notin C$. In the formulation as an integer linear program this translates to an equation $x_{i,j} = 0$ (for $(i + \mathbb{Z}n, j + \mathbb{Z}n) = z$). If there exists an index h , three elements $z_1, z_2, z_3 \in \mathbb{Z}_n^2$, and an automorphism α such that

$$\alpha(\{z_1, z_2, a_1, \dots\}) = \{b_h, z_3, a_1, \dots\}$$

then we can conclude w.l.o.g. that it is not the case that both of z_1 and z_2 are contained in C . In the ILP formulation this translates to an inequality $x_{i,j} + x_{i',j'} \leq 1$ (for $(i + \mathbb{Z}n, j + \mathbb{Z}n) = z_1$ and $(i' + \mathbb{Z}n, j' + \mathbb{Z}n) = z_2$).

With this tool at hand we are able to determine some further exact values of $m_2(\mathbb{Z}_n^2)$ via case differentiations and the **ILOG CPLEX** solver, see Table 2.

In each case we consider ILP (3) augmented by the inequalities arising from the a_i, b_i as described above and by the inequality $\sum_{i=1}^n \sum_{j=1}^n x_{i,j} \geq l + 1$,

where l is the cardinality of a cap C in \mathbb{Z}_n^2 . Since it is not hard to find a cap of cardinality $m_2(\mathbb{Z}_n^2)$ (one can use, e. g., an ILP solver), we can choose $l = m_2(\mathbb{Z}_n^2)$. Thus *resolved* means that the ILP solver has proven that no integer solution can exist in each of the stated subcases. In Table 3 we give the proven exact values of $m_2(\mathbb{Z}_n^2)$ and some bounds which can be obtained by applying the described methods.

n	14	15	16	18	20	21	22	24
$m_2(\mathbb{Z}_n^2)$	12	15	14	17	18	18	18–24	18–24

Table 3. Solving ILP (3) utilizing case differentiations and symmetry

The starting point of our studies was the determination of $m_2(\mathbb{Z}_{25}^2)$. It is not too hard to find a cap of cardinality 20 in \mathbb{Z}_{25}^2 , see Figure 3 for an example. In the projective case very recently, somewhat surprisingly, a cap of cardinality 21 was found [16]. Shortly afterwards Kohnert et al., see [17], verified this constructive result by prescribing a cyclic group of order 3 as a subgroup of the automorphism group of a cap. For the general method of prescribing automorphisms we refer e. g. to [15, 18].

We note that the known bounds for the maximum size of a cap in the projective Hjelmslev plane $\text{PHG}(\mathbb{Z}_{25}^3)$ were 20...25, see e. g. [10]. Here we

conjecture 21 to be the correct value. It is very likely that the method of M. Koch can be continued to an exhaustive search to resolve this case.

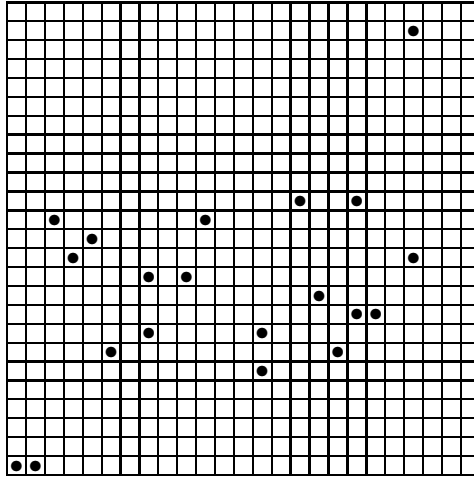


Fig. 3. A cap of cardinality 20 over \mathbb{Z}_{25}^2

To determine $m_2(\mathbb{Z}_{25}^2) = 20$ using the ILP based approach one has to perform several case differentiations, use some geometric insights and exhaustive enumeration. At first we look at pairs $a, b \in \mathbb{Z}_{25}^2$ of points which lie in the same neighborhood \sim_1 , i. e. where a and b are incident with 5^1 common lines, see Section 2. In \mathbb{Z}_{25}^2 two points $(a_1, a_2), (b_1, b_2)$ belong to such a class if and only if both $\hat{a}_1 - \hat{b}_1$ and $\hat{a}_2 - \hat{b}_2$ are divisible by 5. If three pairwise different points $c_1, c_2, c_3 \in \mathbb{Z}_{25}^2$ fulfill $a \sim_1 b \sim_1 c$ then they are collinear. Thus in a cap no such triple can exist.

Now we consider caps C which contain at least three pairs $c_1 \sim_1 c_2, c_3 \sim_1 c_4$, and $c_5 \sim_1 c_6$. Since C is a cap we have $c_1 \not\sim_1 c_3, c_1 \not\sim_1 c_5$, and $c_3 \not\sim_1 c_5$. Due to the symmetry group G , defined above as a subgroup of the automorphism group, we can assume w.l.o.g. $c_1 = (\bar{0}, \bar{0}), c_3 = (\bar{1}, \bar{0}),$ and $c_5 = (\bar{0}, \bar{1})$. For $\{c_1, \dots, c_6\}^G$ there are 104 orbits. In each of these 104 cases the augmented ILP (3) has an optimal solution less than 21.

Thus we only have to consider caps with at most two pairs $c_1 \sim_1 c_2$ and $c_3 \sim_1 c_4$. So in the next step we consider all 18 orbits $\{c_1, \dots, c_5\}^G$ with $c_1 = (\bar{0}, \bar{0}), c_3 = (\bar{1}, \bar{0}), c_1 \sim_1 c_4,$ and $c_2 \sim_1 c_5$. As additional inequalities we can use that no further pair $c_i \sim_1 c_j$ can exist in C . Also, here the optimal solution of the augmented ILPs (3) is less than 21.

In a further step we consider all 39 orbits $\{c_1, \dots, c_4\}^G$ with $c_1 = (\bar{0}, \bar{0})$, $c_3 = (\bar{1}, \bar{0})$, and $c_1 \sim_1 c_4$. Again the optimal solution of the augmented ILPs (3) is less than 21. Thus in a possible cap of cardinality at least 21 in \mathbb{Z}_{25}^2 no pair $c_i \sim_1 c_j$ can exist.

Next we look at the orbits of $\{c_1, \dots, c_4\}^G$ with $c_1 = (\bar{0}, \bar{0})$, $c_3 = (\bar{1}, \bar{0})$ where no pair $c_i \sim_1 c_j$ occurs. We pick a representative of each orbit and label them by z_1, \dots, z_{33} . For $i \geq 4$ we set $a_1 = c_1$, $a_2 = c_2$, $a_3 = c_3$, $a_4 = z_i$, $b_1 = z_1, \dots, b_{i-1} = z_{i-1}$ in order to prescribe or to forbid some points of cap C . In each of these cases we can verify that C must contain less than 21 elements. Thus only three possibilities for $\{c_1, \dots, c_4\}^G$ are left.

For the final step we utilize orderly generation [22] with some look ahead: we run through all 25 equivalence classes $\{x \in \mathbb{Z}_{25} \mid x \sim_1 y\}$. In each equivalence class we can pick at most one point. So either we pick one element of such an equivalence class or we decide not to take an element of this equivalence class. If we have fixed some elements of our cap C then it may happen that there are equivalence classes where we have no possibility to pick an element due to the restriction of at most 2 points on a line. So for each partial cap we can count the number r of equivalence classes, where it is possible to select a further point for cap C . If $|C| + r < 21$ then we can stop extending cap C .

4.2. Minimum cardinality of complete caps over \mathbb{Z}_n^2 . From [13] we can cite the bounds

$$\max\left(4, \sqrt{2p} + \frac{1}{2}\right) \leq n_2(\mathbb{Z}_n^2) \leq \max(4, p + 1),$$

where p is the smallest prime divisor of n , and

$$n_2(\mathbb{Z}_{nm}^2) \leq \min\left(n_2(\mathbb{Z}_n^2), n_2(\mathbb{Z}_m^2)\right)$$

for coprime integers n and m . Additionally the author of [13] conjectures $n_2(\mathbb{Z}_{p^a}^2) \leq n_2(\mathbb{Z}_{p^b}^2)$ for $a \leq b$, p being a prime and $n_2(\mathbb{Z}_{nm}^2) = \min\left(n_2(\mathbb{Z}_n^2), n_2(\mathbb{Z}_m^2)\right)$ for coprime integers n and m . If these conjectures turn out to be true then it would suffice to determine the values $n_2(\mathbb{Z}_p^2)$ for primes p .

If n is divisible by 2 or 3 then we can conclude $n_2(\mathbb{Z}_n^2) = 4$ from the above inequalities and $n_2(\mathbb{Z}_2^2) = n_2(\mathbb{Z}_3^2) = 4$. For $n > 1$ we have $n_2(\mathbb{Z}_n^2) \geq 4$. So in Table and Table we have given the exact values and bounds for $n_2(\mathbb{Z}_n^2)$ arising from ILP (3), where n is either a prime or coprime to 6. We would like to remark that some of these numbers are already given in [13]. If n is a prime then we can

assume w.l.o.g. that the points $(\bar{0}, \bar{0})$, $(\bar{1}, \bar{0})$, and $(\bar{0}, \bar{1})$ are contained in cap C . In this case some results from [7] can be used, e. g. we have $n_2(\mathbb{Z}_n^2) \geq t(2, q) - 2$, where $t(2, q)$ denotes the smallest size of a complete cap in $PG(2, q)$. For the exact values of $t(2, q)$ for $q \leq 29$ we refer to [9, 20]. In [7] there is mentioned a construction that produces a cap of size $\frac{q-3}{2}$ which is complete at least for $q > 413$. The exact value of $n_2(\mathbb{Z}_{25}^2)$ is known to be 6 but unfortunately we were not able to validate the lower bound using our ILP based approach.

n	2	3	5	7	11	13	17	19	23	25	29
$n_2(\mathbb{Z}_n^2)$	4	4	5	6	7	8	8...10	8...11	8...12	4...6	11...15

Table 4. Values of $n_2(\mathbb{Z}_n^2)$ for small n which are either prime or coprime to 6

n	31	37	41	43	47
$n_2(\mathbb{Z}_n^2)$	9...16	10...17	10...20	10...21	11...22

Table 5. Values of $n_2(\mathbb{Z}_n^2)$ for small n which are either prime or coprime to 6

4.3. Maximum cardinality of caps over \mathbb{Z}_n^2 which are subsets of permutations. Obviously we have $\sigma(\mathbb{Z}_n^2) \leq n$ since a permutation consists of n points. Utilizing the ILP formulation (4) and the ILOG CPLEX solver we have obtained the values and bounds of Table and Table . So e. g. for $n \in \{1, 2, 4, 6, 8, 12\}$ there exist permutations whose graphs in \mathbb{Z}_n^2 are caps. We would like to remark that for this problem the applicable group of automorphisms is much smaller than for the other two problems. Here translations, changes of the coordinate axes, and reflection at one of the axes are automorphisms.

n	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19
$\sigma(n)$	1	2	2	4	4	6	6	8	6	8	10	12	12	12	13	13	16	13	18

Table 6. Values and bounds for $\sigma(\mathbb{Z}_n^2)$

n	20	21	22	23	24	25	26	27	28	29	30
$\sigma(n)$	16	16-17	16-17	22	20-22	19-22	18-24	18-25	22-27	28	22-29

Table 7. Values and bounds for $\sigma(\mathbb{Z}_n^2)$

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Received March 5, 2009
Final Accepted May 15, 2009