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### Spanning trails containing given edges

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#### Abstract

A graph G is Eulerian-connected if for any u and v in V(G), G has a spanning (u, v)-trail. A graph G is edge-Eulerian-connected if for any e' and e'' in E(G), G has a spanning (e', e'')-trail. For an integer  $r \ge 0$ , a graph is called r-Eulerian-connected if for any  $X \subseteq E(G)$  with  $|X| \le r$ , and for any  $u, v \in V(G)$ , G has a spanning (u, v)-trail T such that  $X \subseteq E(T)$ . The r-edge-Eulerianconnectivity of a graph can be defined similarly. Let  $\theta(r)$  be the minimum value of k such that every k-edge-connected graph is r-Eulerian-connectivity are also discussed.

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#### 1. Introduction

We follow the notation of Bondy and Murty [1], except that graphs have no loops. A graph G is Hamiltonianconnected if for every pair of vertices u, v of G, there is a Hamiltonian (u, v)-path in G. For a graph G, a trail is a vertex-edge alternating sequence  $v_0, e_1, v_1, e_2, \ldots, e_{k-1}, v_{k-1}, e_k, v_k$  such that all the  $e_i$ 's are distinct and  $e_i = v_{i-1}v_i$ for all *i*. Let  $e', e'' \in E(G)$ . A trail in G whose first edge is e' and whose last edge is e'' is called an (e', e'')-trail. For  $u, v \in V(G)$ , a (u, v)-trail of G is a trail in G whose origin is u and whose terminus is v. A trail H is called a dominating trail of G if every edge of G is incident with at least one vertex of H in G. A trail H is called a spanning trail if V(H) = V(G). If u = v, then a (u, v)-trail in G is a closed trail, which is also called a *Eulerian subgraph* of G. A graph is called supereulerian if it has a spanning closed trail. The collection of all supereulerian graphs is denoted by  $\mathscr{SL}$ .

A graph G is Eulerian-connected if for any u, v in V(G) (including the case u = v), G has a spanning (u, v)-trail. A graph is called r-Eulerian-connected if for any  $X \subseteq E(G)$  with  $|X| \leq r$ , and for any  $u, v \in V(G)$ , G has a spanning (u, v)-trail *T* such that  $X \subseteq E(T)$ . For an integer  $r \ge 0$ , the collection of all *r*-Eulerian-connected graphs is denoted by  $\mathscr{EL}(r)$ . Obviously,  $\mathscr{EL}(r) \subseteq \mathscr{SL}$  for all  $r \ge 0$ .

A graph *G* is *edge-Eulerian-connected* if for any e', e'' in E(G), *G* has a spanning (e', e'')-trail. A graph is called *r*-edge-Eulerian-connected if for any  $X \subseteq E(G)$  with  $|X| \leq r$  and for any e',  $e'' \in E(G)$ , *G* has a spanning (e', e'')-trail *T* such that  $X \subseteq E(T)$ . For an integer  $r \geq 0$ , the collection of all *r*-edge-Eulerian-connected graphs is denoted by  $\mathscr{E}(r)$ .

Many studies have been done on Eulerian graphs (see [7]). For the literature on the subject of supereulerian graphs, see surveys [3,6]. Harary and Nash-Williams [9] demonstrated the relationship between Eulerian subgraphs and Hamiltonian cycles in the line graph of *G*. Zhan [14] studied (e', e'')-trails of a graph *G* for the Hamiltonian connectivity of the line graph of *G*. In the study of spanning trails of graphs [2], Catlin introduced the concept of *collapsible graphs*. For a graph *G*, let O(*G*) be the set of odd degree vertices of *G* and let *R* be an even subset of V(G). A subgraph  $H_R$  of *G* is called a *spanning R-trail* if  $H_R$  is a spanning connected subgraph such that  $O(H_R) = R$ . A graph *G* is *collapsible* if for every even subset  $R \subseteq V(G)$ , *G* has a spanning *R*-trail. We will regard an empty set as an even subset and  $K_1$  as both collapsible and supereulerian. The collection of all collapsible graphs is denoted by  $\mathscr{CL}$ . By the definition of collapsible graphs, we have:

#### **Proposition A.** Let G be a collapsible graph. Then each of the following holds

- (i) *G* is supereulerian.
- (ii) G is Eulerian-connected.

**Proof.** For any vertices  $u, v \in V(G)$ . Let  $R = \emptyset$  if u = v, or  $R = \{u, v\}$  if  $u \neq v$ . Since G is collapsible, it has a spanning subgraph  $H_R$  such that  $O(H_R) = R$ . Therefore,  $H_R$  is a spanning Eulerian subgraph of G if  $R = \emptyset$ , or  $H_R$  is a (u, v)-spanning trail of G.  $\Box$ 

Let  $X \subseteq E(G)$  and let R be an even subset of V(G). A spanning R-trail  $H_R$  of G such that  $X \subseteq E(H_R)$  is called *a* spanning (R, X)-trail, and denoted by  $H_R(X)$ . A graph is called strongly r-Eulerian-connected if for any  $X \subseteq E(G)$  with  $|X| \leq r$  and for any even subset  $R \subseteq V(G)$ , G has a spanning R-trail  $H_R$  such that  $X \subseteq E(H_R)$  (i.e. G has a  $H_R(X)$ ). The collection of all strongly r-Eulerian-connected graphs is denoted by  $\mathscr{GE}(r)$ .

For an integer *r*, define  $\mathscr{L}(r)$  to be the family of graphs such that  $G \in \mathscr{L}(r)$  if and only if for any subset  $X \subseteq E(G)$ with  $|X| \leq r$ , *G* has an spanning Eulerian subgraph *H* such that  $X \subseteq E(H)$ . Define f(r) to be the minimum value of *k* such that every *k*-edge-connected graph *G* is in  $\mathscr{L}(r)$ . In [12], Lai found f(r) for all the values of *r* (see Corollary 3.6). Let  $\theta(r)$  be the minimum value of *k* such that every *k*-edge-connected graph is in  $\mathscr{E}\mathscr{L}(r)$  and let  $\psi(r)$  be the minimum value of *k* such that every *k*-edge-connected graph is in  $\mathscr{S}\mathscr{E}(r)$ . Since  $\mathscr{S}\mathscr{E}(r) \subseteq \mathscr{E}\mathscr{L}(r) \subseteq \mathscr{L}(r)$ ,

$$f(r) \leqslant \theta(r) \leqslant \psi(r). \tag{1}$$

Let  $\xi(r)$  be the minimum value of k such that every k-edge-connected graph is in  $\mathscr{E}\mathscr{E}(r)$ . In this paper, we will determine the values of  $\theta(r)$ ,  $\psi(r)$ , and  $\xi(r)$  for all  $r \ge 0$ .

In the next section, we will present Catlin's reduction method and some preliminary results which are needed in our proofs. Our main results are in Sections 3 and 4. We will present our results on *r*-Eulerian-connected graphs, and give the values of  $\theta(r)$  and  $\psi(r)$  for all  $r \ge 0$ . Section 4 contains results on the *r*-edge-Eulerian connected graphs.

#### 2. Catlin's reduction method and preliminary results

Let *H* be a connected subgraph of *G*. The contraction G/H is obtained from *G* by contracting each edge of *H* and deleting the resulting loops. In [2], Catlin showed that every graph *G* has a unique collection of pairwise vertex-disjoint maximal collapsible subgraphs  $H_1, H_2, \ldots, H_k$  such that  $\bigcup_{i=1}^k V(H_i) = V(G)$ . The *reduction* of *G* is obtained from *G* by contracting each of  $H_i$  into a vertex  $v_i$  for all *i*, and is denoted by *G'*. Each  $H_i$  is called a preimage of  $v_i$  in *G*, and  $v_i$  is called the contraction image of  $H_i$  in *G'*. A vertex *v* in *G'* is called a *trivial contraction* if its preimage in *G* is  $K_1$ . A graph *G* is reduced if *G* is the reduction of some graph. Let F(G) be the minimum number of edges that must be added to *G* so that the resulting graph has 2 edge-disjoint spanning trees.

**Theorem 2.1** (*Catlin* [2]). Let G be a graph, and let G' be the reduction of G. Each of the following holds.

- (i) G is supereulerian if and only if G' is supereulerian.
- (ii) *G* is collapsible if and only if  $G' \cong K_1$
- (iii) |E(G')| + F(G') = 2|V(G')| 2.

In [10], Jaeger proved that a graph with two edge-disjoint spanning trees is supereulerian. In [2], Catlin proved that if G has two edge-disjoint spanning trees, then G is collapsible. It is well known now that a 2k-edge-connected graph has k edge-disjoint spanning trees [8,11,13]. Thus, we have:

**Theorem 2.2.** If G is 4-edge-connected, then G is collapsible.

In [4], Catlin proved:

**Theorem 2.3** (*Catlin* [4]). Let G be a graph and let  $k \ge 1$  be an integer. The following are equivalent:

- (i) *G* is 2*k*-edge-connected;
- (ii) For any  $X \subseteq E(G)$  with  $|X| \leq k$ , G X has k edge-disjoint spanning trees.

**Corollary 2.4** (*Catlin* [4]). Let G be a graph and let  $k \ge 1$  be an integer. The following are equivalent:

- (i) G is (2k + 1)-edge-connected;
- (ii) For any  $X \subseteq E(G)$  with  $|X| \leq k + 1$ , G X has k-edge-disjoint spanning trees.

The following theorems will be needed in our proofs.

**Theorem 2.5** (*Catlin et al.* [5]). Let G be a connected graph. If  $F(G) \leq 2$ , then either G is collapsible, or the reduction of G is in  $\{K_2, K_{2,t} : t \geq 1\}$ .

Let *e* be an edge in *G*. Edge *e* is *subdivided* when it is replaced by a path of length 2 whose internal vertex, denoted by v(e), has degree 2 in the resulting graph. The process of taking an edge *e* and replacing it by that path of length 2 is called subdividing *e*. Let *G* be a graph and let  $X \subseteq E(G)$ . Let  $G_X$  be the graph obtained from *G* by subdividing each edge in *X*. Then  $V(G_X) = V(G) \cup \{v(e) \text{ for each } e \in X\}$ .

**Lemma 2.6.** Let  $k \ge 2$  be an integer. Let *G* be a connected graph and let  $X \subseteq E(G)$ . Let *R* be an even subset of V(G). Then each of the following holds

- (i) G has a spanning (R, X)-trail  $H_R(X)$  if and only if  $G_X$  has a spanning R-trail. In particular, G has a spanning closed trail H such that  $X \subseteq E(H)$  if and only if  $G_X$  is supereulerian.
- (ii) If  $G_X$  is collapsible, then  $G_X$  has a spanning *R*-trail.
- (iii) Let  $X = X_1 \cup X_2$  with  $X_1 \cap X_2 = \emptyset$ . Then  $F(G_X) \leq F((G X_1)_{X_2})$ .
- (iv) If G has k edge-disjoint spanning trees, then for any  $X \subseteq E(G)$  with  $|X| \leq 2k 2$ ,  $F(G_X) \leq 2$ .

**Proof.** (i) and (ii) follow from the definitions of collapsibility and  $G_X$ .

(iii) Let  $p = F((G - X_1)_{X_2})$ . Let  $E_p$  be the p edge set such that  $(G - X_1)_{X_2} + E_p$  has 2-edge-disjoint spanning trees  $(T_1 \text{ and } T_2)$ . Let  $X_1 = \{e_1, e_2, \dots, e_s\}$  and each  $e_i = u_i v_i$   $(1 \le i \le s)$ . By the definition of  $G_X$ , we know that  $G_X$  can be obtained from  $(G - X_1)_{X_2}$  by joining each pair of  $u_i$  and  $v_i$  by a path  $P_i = u_i v(e_i)v_i$  where  $v(e_i)$  is a new vertex. Therefore,  $T_1 + \bigcup_{i=1}^s \{u_i v(e_i)\}$  and  $T_2 + \bigcup_{i=1}^s \{v(e_i)v_i\}$  are two edge-disjoint spanning trees in  $G_X + E_p$ , and so  $F(G_X) \le p = F((G - X_1)_{X_2})$ .

(iv) Let  $T_1, T_2, \ldots, T_k$  be k edge-disjoint spanning trees of G. Without lost of generality, we may assume that

$$|X \cap E(T_1)| \leq |X \cap E(T_2)| \leq \dots \leq |X \cap E(T_k)|.$$
<sup>(2)</sup>

Since  $k \ge 2$ ,  $|X| \le 2k - 2$ ,  $T_i$ 's are edge-disjoint, and by (2),

 $|X \cap E(T_1)| + |X \cap E(T_2)| \leq 2.$ 

Let  $X = \{e_1, e_2, \dots, e_p\}$  where  $p \leq 2k - 2$ , and let  $e_i = u_i v_i$  for all  $1 \leq i \leq p$ . Since  $G_X$  is the graph obtained from G by subdividing  $e_i$   $(1 \leq i \leq p)$ ,  $V(G_X) = V(G) \cup \{v(e_i) : 1 \leq i \leq p\}$ , and  $E(G_X) = (E(G) - X) \cup \{u_i v(e_i), v(e_i)v_i : 1 \leq i \leq p\}$ .

*Case* 1.  $|X \cap E(T_1)| + |X \cap E(T_2)| = 0$ .

Then  $T_1 + \bigcup_{i=1}^p \{u_i v(e_i)\}$  and  $T_2 + \bigcup_{i=1}^p \{v(e_i)v_i\}$  are two edge-disjoint spanning trees in  $G_X$  and so  $F(G_X) = 0 \leq 2$ . Case 2.  $|X \cap E(T_1)| + |X \cap E(T_2)| = 1$ .

By (2) and (3),  $|X \cap E(T_1)| = 0$  and  $|X \cap E(T_2)| = 1$ . Let  $e_2 = u_2v_2$  be the edge in  $X \cap E(T_2)$ . Then  $T'_2 = T_2 - e_2 + \{u_2v(e_2), v(e_2)v_2\} \bigcup_{i\neq 2}^p \{v(e_i)v_i\}$  is a spanning tree in  $G_X$ . To obtain another spanning tree which covers  $v(e_2)$ , we can add an edge  $e' = u_1v(e_2)$  to  $G_X$ . Then  $T'_1 = T_1 + \{e'\} \bigcup_{i\neq 2}^p \{u_iv(e_i)\}$  is a spanning tree in  $G_X + e'$ . Therefore,  $T'_1$  and  $T'_2$  are two edge-disjoint spanning trees in  $G_X + e'$ . This shows that  $F(G_X) = 1 \leq 2$ .

Case 3.  $|X \cap E(T_1)| + |X \cap E(T_2)| = 2.$ 

By (2) and (3), either  $|X \cap E(T_1)| = |X \cap E(T_2)| = 1$ , or  $|X \cap E(T_1)| = 0$  and  $|X \cap E(T_2)| = 2$ . We prove  $F(G_X) \le 2$  for the case  $|X \cap E(T_1)| = |X \cap E(T_2)| = 1$  here. The case  $|X \cap E(T_1)| = 0$  and  $|X \cap E(T_2)| = 2$  can be proved similarly.

Let  $e_1 \in X \cap E(T_1)$  and  $e_2 \in X \cap E(T_2)$ . Then  $T'_1 = T_1 - e_1 + \{u_1v(e_1), v(e_1)v_1\} \bigcup_{i=3}^p \{u_iv(e_i)\}$  is a tree containing  $V(G_X) - v(e_2)$ , and  $T'_2 = T_2 - e_2 + \{u_2v(e_2), v(e_2)v_2\} \bigcup_{i=3}^p \{v(e_i)v_i\}$  is a tree containing  $V(G_X) - v(e_1)$ . Therefore, adding two new edges  $e' = u_1v(e_2)$  and  $e'' = v(e_1)v_2$  to  $G_X$ , we have two edge-disjoint spanning trees  $T'_1 + e'$  and  $T'_2 + e''$  in  $G_X + \{e', e''\}$ . This shows that  $F(G_X) \leq 2$ . The proof is complete.  $\Box$ 

**Lemma 2.7.** Let G be a graph with  $\kappa'(G) \ge 3$ , and let  $X \subseteq E(G)$ . Let  $G_X$  be the graph obtained from G by subdividing each edge in X. If the reduction of  $G_X$  is  $K_{2,t}$ , then each of the following holds.

- (i) Every degree 2 vertex in  $G'_X$  is a vertex obtained by subdividing an edge in X.
- (ii)  $|X| \ge t \ge \kappa'(G)$ , and X is an edge cut of G.
- (iii) There is a subset  $X_1 \subseteq X$  with  $t = |X_1|$  such that each path between the two vertices of degree t in  $K_{2,t}$  is obtained by subdividing an edge in  $X_1$ . Furthermore,  $G_X - X_1$  has only two collapsible components (say  $H_1$  and  $H_2$ ) such that  $V(G_X) = V(H_1) \cup V(H_2) \bigcup_{e \in E_1} \{v(e)\}$ , and  $G'_X = K_{2,t}$  is obtained by contracting  $H_1$  and  $H_2$  (i.e.  $G'_X = (G_X/H_1)/H_2 = K_{2,t}$ ).

**Proof.** Let  $E(G'_X) = E(K_{2,t}) = \{uw_i, w_iv\}$   $(1 \le i \le t)$  where each  $w_i$  is a degree 2 vertex in  $G'_X$ . Note that  $w_i$  is a trivial contraction, and (i) holds. Otherwise the two edges incident with  $w_i$  will form an edge-cut of G, contrary to that  $\kappa'(G) \ge 3$ . Hence, each path  $uw_iv$  is obtained by subdividing an edge in X and so  $t \le |X|$ .

Let  $E' = \{uw_i : 1 \le i \le t\}$ . Then E' is an edge-cut of  $G'_X$ . Since each path  $uw_iv$  in  $G_X$  is obtained by subdividing an edge  $e \in X \subseteq E(G)$ , we have an edge set  $X_1 \subseteq X$  such that each edge in  $X_1$  corresponding to a path  $uw_iv$  in  $G_X$ , and  $|X_1| = |E'| = t$ . Therefore,  $X_1$  is an edge cut in G. Since  $X_1 \subseteq X$ , X is an edge-cut of G and  $|X| \ge |E'| = t \ge \kappa'(G)$ .

Note  $V(G'_X) = \{u, v, w_i : 1 \le i \le t\}$  where d(u) = d(v) = t. Let  $H_1$  be the preimage of u, and let  $H_2$  be the preimage of v. Therefore,  $G'_X$  is obtained by subdividing each edge in  $X_1$ , and then contracting  $H_1$  and  $H_2$ , respectively. Statement (iii) is proved.  $\Box$ 

**Lemma 2.8.** Let G be an r-edge-connected graph  $(r \ge 4)$ . Let  $X \subseteq E(G)$ . Let  $G_X$  be the graph obtained from G by subdividing each edge in X. Let  $G'_X$  be the reduction of  $G_X$  and let  $V_r$  be the set of vertices of degree less than r in  $G'_X$ . Let  $D_i = \{v \in V(G'_X) : d(v) = i\}$   $(i \ge 2)$ . If  $F(G'_X) \ge 3$ , then each of the following holds:

(i) each vertex in  $V_r$  has degree 2 (i.e.  $V_r = D_2$ ) and  $|V_r| \leq |X|$ .

(ii)  $(r-4)|V(G'_X)| + 10 \leq (r-2)|V_r| \leq (r-2)|X|$ .

(iii)  $10 + (r-4)|D_r| + (r-3)|D_{r+1}| + \dots + \leq 2|V_r| \leq 2|X|.$ 

**Proof.** Since the degree of each vertex u in  $V_r$  is less than r, u must be a trivial contraction in  $G'_X$ . Otherwise, the edges incident with u will form an edge cut with size less than r, contrary to  $\kappa'(G) \ge r$ . Therefore,  $V_r \subseteq V(G_X) - V(G)$ ,

a subset of the vertices obtained in the process of subdividing each edge in X. Thus each vertex in  $V_r$  has degree 2 and

$$|V_r| \leqslant |X|. \tag{4}$$

Let  $c = |V(G'_X)|$ . Since  $F(G'_X) \ge 3$ , by (iii) of Theorem 2.1,

$$|E(G'_X)| = 2|V(G'_X)| - 2 - F(G'_X) \leq 2c - 5.$$

Hence,

$$\sum_{v \in V(G'_X)} d(v) = 2|E(G'_X)| \leqslant 4c - 10.$$
(5)

Since  $\kappa'(G_X) \ge 2$ ,  $\delta(G'_X) \ge 2$ . Then by (5)

$$2|V_r| + r(c - |V_r|) \leq 2|V_r| + \sum_{v \notin V_r} d(v) = \sum_{v \in V(G'_X)} d(v) = 2|E(G'_X)| \leq 4c - 10.$$
(6)

By (4), (6), and  $c = |V(G'_X)|$ ,

$$(r-4)|V(G'_X)| + 10 \leqslant (r-2)|V_r| \leqslant (r-2)|X|.$$
(7)

By (6), and  $V(G'_X) = V_r \bigcup_{i=r} D_i$ ,

$$2|V_r| + r|D_r| + (r+1)|D_{r+1}| + \dots \leq 4(|V_r| + |D_r| + |D_{r+1}| + \dots) - 10.$$

Hence,

$$10 + (r-4)|D_r| + (r-3)|D_{r+1}| + \dots \leq 2|V_r| \leq 2|X|.$$

**Lemma 2.9.** Let G be a graph and let  $e_1, e_2 \in E(G)$  and let  $X \subseteq E(G)$ . Let  $X_0 = X \cup \{e_1, e_2\}$ . Let  $G_{X_0}$  be the graph obtained from G by subdividing each edge in  $X_0$ . Let  $v(e_1)$  and  $v(e_2)$  be the two vertices subdividing  $e_1$  and  $e_2$ , respectively. Then

(i) If  $G_{X_0}$  has a spanning  $(v(e_1), v(e_2))$ -trail, then G has a spanning  $(e_1, e_2)$ -trail containing X.

(ii) If  $G_{X_0}$  is collapsible, then G has a spanning  $(e_1, e_2)$ -trail containing X.

**Proof.** Follows from the definitions of collapsibility and  $G_{X_0}$ .

#### 3. The *r*-Eulerian-connected graphs

The Petersen graph and many other 3-edge-connected graphs have no spanning closed trails. Thus, for any  $r \ge 0$ ,  $\psi(r) \ge \theta(r) \ge 4$ . By Theorem 2.2, we know that  $\psi(0) = \theta(0) = 4$ . The following example shows that for  $r \ge 3$ ,  $\psi(r) \ge \theta(r) \ge r + 1$ .

**Example 1.** Let  $r \ge 3$  be an integer, and let n and m be two integers such that  $n \ge r + 1$  and  $m \ge r + 1$ . Let  $G_1 = K_n$  with  $V(G_1) = \{u_1, u_2, \ldots, u_n\}$ , and let  $G_2 = K_m$  with  $V(G_2) = \{v_1, v_2, \ldots, v_m\}$ . Define the graph G to be the graph obtained from  $G_1$  and  $G_2$  by connecting  $G_1$  and  $G_2$  with the new edge set  $X = \{e_1, e_2, \ldots, e_r\}$  where  $e_i = u_i v_i$  for all  $i = 1, 2, \ldots, r$ . Then G is an r-edge-connected graph. If r is even, then we choose u from  $G_1$ , and v from  $G_2$ . If r is an odd integer, then we choose u and v both from  $G_1$ . Then G has no spanning (u, v)-trails containing all the edges of X. This example also shows that G has no spanning (e', e'')-trails containing all the edges of X for some pair of  $e', e'' \in E(G)$ . See Fig. 1 below for the case r = 4 where  $X = \{e_1, e_2, e_3, e_4\}$  and  $G_1 \cong G_2 \cong K_5$ . This shows that  $\psi(r) \ge \theta(r) \ge r + 1$ . In the following, we will show that  $\psi(r) = \theta(r) = r + 1$ .

This example suggests the following necessary condition for *r* Eulerian-connected graphs, and the lower bounds for  $\psi(r)$ ,  $\theta(r)$  and  $\xi(r)$ .



Fig. 1.

**Theorem 3.0.** Let  $r \ge 3$ . Then  $\psi(r) \ge \theta(r) \ge r + 1$  and  $\xi(r) \ge r + 1$ . Furthermore, if G is an r-Eulerian-connected graph, then G is (r + 1)-edge-connected.

**Proof.** By way of contradiction, suppose that the edge-connectivity of *G* is  $k \le r$ . Let *X* be an edge cut with |X| = k and let  $H_1$  and  $H_2$  be two components of G - X. If |X| = k is even, we can choose a vertex *u* from  $H_1$  and a vertex *v* from  $H_2$ . Then *G* has no spanning (u, v) trail that contains *X*, a contradiction. If |X| = k is odd, then we can choose a vertex *u* from  $H_1$ . Since *X* has odd number of edges, *G* does not have a closed trail that starts and ends at *u* containing *X*, a contradiction again.  $\Box$ 

For a real number x, let  $\lfloor x \rfloor$  be the largest integer that is less than or equal to x.

**Theorem 3.1.** Let  $r \ge 4$  be an integer and let  $k = \lfloor \frac{r}{2} \rfloor$ . Let *G* be an *r*-edge-connected graph and let  $X \subseteq E(G)$  with  $|X| \le r + k - 2$ . it Then one of the following holds:

(i)  $G_X$  is collapsible, or

(ii) *X* is an edge cut of *G* and  $|X| \ge r$ .

**Proof.** Let  $X \subseteq E(G)$  with  $|X| \leq r + k - 2$ . Define  $G_X$  as before and assume that  $G_X$  is not collapsible. We will show that the reduction  $G'_X$  is  $K_{2,t}$  with  $t \geq 2$  first. Consider the following two cases:

*Case* 1. *r* is even. Then r = 2k, and  $|X| \leq 3k - 2$ .

Since  $|X| \leq 3k - 2$ , we can choose a subset  $X_1$  of X and let  $X_2 = X - X_1$ , such that  $|X_1| \leq k$  and  $|X_2| \leq 2k - 2$ . By Theorem 2.3,  $G - X_1$  has k-edge-disjointed spanning trees. Then by Lemma 2.6(iv),  $F((G - X_1)_{X_2}) \leq 2$ . By Lemma 2.6(iii),  $F(G_X) \leq F((G - X_1)_{X_2}) \leq 2$ . Since  $G_X$  is not collapsible, by Theorem 2.5,  $G'_X \in \{K_2, K_{2,t}\}$   $(t \geq 1)$ . Since G is r-edge-connected  $(r \geq 4)$ ,  $G_X$  is 2-edge-connected. Therefore,  $G'_X = K_{2,t}$   $(t \geq 2)$ .

Case 2. r is odd. Then r = 2k + 1 and  $|X| \leq 3k - 1$ .

Let  $X_1$  be a subset of X and let  $X_2 = X - X_1$  such that  $|X_1| \le k + 1$  and  $|X_2| \le 2k - 2$ . By Corollary 2.4,  $G - X_1$  has k-edge-disjointed spanning trees. By Lemma 2.6(iii) and (iv),  $F(G_X) \le F((G - X_1)_{X_2}) \le 2$ . Using the same argument for the case 1 above, we have  $G'_X = K_{2,t}$  ( $t \ge 2$ ).

Therefore, by Lemma 2.7, Theorem 3.1 is proved.  $\Box$ 

From the proof of Theorem 3.1, we have the following:

**Theorem 3.1'.** Let  $r \ge 4$  be an integer and let  $k = \lfloor \frac{r}{2} \rfloor$ . Let *G* be an *r*-edge-connected graph. Let  $X \subseteq E(G)$  with  $|X| \le r + k - 2$  and let  $G_X$  be the graph obtained from *G* by subdividing every edge in *X*. Let  $G'_X$  be the reduction of  $G_X$ . Then exactly one of the following holds

- (i)  $G_X$  is collapsible, or
- (ii)  $G_X$  can be contracted to  $K_{2,t}$  (i.e.  $G'_X = K_{2,t}$ ) in such a way that each degree vertex in  $K_{2,t}$  is a trivial contraction and  $r \le t \le |X|$ .

**Theorem 3.2.** Let  $r \ge 4$  be an integer and let  $k = \lfloor \frac{r}{2} \rfloor$ . Let *G* be an *r*-edge-connected graph. Let  $X \subseteq E(G)$  with  $|X| \le r + k - 2$ . Then one of the following holds

- (i) for any even subset  $R \subseteq V(G)$ , G has a spanning R-trail  $H_R$  such that  $X \subseteq E(H_R)$ , or
- (ii) *X* is an edge cut of *G* and  $|X| \ge r$ .

**Proof.** For a given edge set  $X \subseteq E(G)$ , by Lemma 2.6(ii), if  $G_X$  is collapsible, then G has a spanning (R, X)-trail for any even subset  $R \subseteq V(G)$ . Theorem 3.2 follows from Theorem 3.1.  $\Box$ 

**Corollary 3.3.** Let  $r \ge 4$  be an integer, and let  $k = \lfloor \frac{r}{2} \rfloor$ . Let *G* be an *r*-edge-connected graph. Let  $X \subseteq E(G)$  with  $|X| \le r + k - 2$ . If *X* is not an edge cut of *G*, then *G* has a spanning (R, X)-trail for any even subset  $R \subseteq V(G)$ .

**Proof.** Following Theorem 3.1 and Lemma 2.6 immediately.

**Corollary 3.4.** Let  $r \ge 3$ . Then G is strongly r-Eulerian-connected if and only if G is (r + 1)-edge-connected.

**Proof.** The necessary condition follows from Theorem 3.0. For the sufficient condition, let  $X \subseteq E(G)$  with  $|X| \leq r$ . Then  $|X| < \kappa'(G) = r + 1$ . *X* is not an edge cut of *G* and by Theorem 3.2, the statement holds.  $\Box$ 

**Theorem 3.5.** Let  $r \ge 0$ . Then

 $\psi(r) = \theta(r) = \begin{cases} 4 & \text{if } 0 \leq r \leq 2, \\ r+1 & \text{if } r \geq 3. \end{cases}$ 

**Proof.** Since there exist 3-edge-connected graphs which are not supereulerian,  $\psi(r) \ge \theta(r) \ge 4$  for  $r \ge 0$ . By Theorem 3.1, if *G* is 4-edge-connected, then any edge set *X* with  $|X| \le 2$  can not be an edge cut of *G*. Therefore  $G_X$  is collapsible, and so  $\theta(r) = \psi(r) \le 4$  if  $r \le 2$ . For  $r \ge 3$ , it follows from Corollary 3.4 that  $\psi(r) = \theta(r) = r + 1$ .  $\Box$ 

**Corollary 3.6** (*Lai* [12]). Let  $r \ge 0$  be an integer. Then

$$f(r) = \begin{cases} 4, & 0 \leq r \leq 2, \\ r+1, & r \geq 3 \text{ and } r \text{ is odd,} \\ r, & r \geq 4 \text{ and } r \text{ is even.} \end{cases}$$

**Proof.** Since there exist 3-edge-connected graphs that are not supereulerian,  $f(r) \ge 4$ . Since  $f(r) \le \theta(r)$ , by Theorem 3.1, f(r) = 4 if  $r \le 2$ . For  $r \ge 3$ , if r is odd, Example 1 with an odd number r shows that  $f(r) \ge r + 1$ . By Theorem 3.1, since  $f(r) \le \theta(r) \le r + 1$ , f(r) = r + 1 if r is odd. If r is even, by Theorem 3.1', for any r-edge-connected graph G and any  $X \subseteq E(G)$  with  $|X| \le r$ , either  $G_X$  is collapsible or the reduction  $G'_X \cong K_{2,r}$ . Since  $K_{2,r}$  is supereulerian when r is even and all collapsible graphs are supereulerian,  $G_X$  is supereulerian. Then by Lemma 2.6(i), G has a spanning Eulerian subgraph H with  $X \subseteq E(H)$ . Therefore, f(r) = r if r is even.  $\Box$ 

Corollary 3.6 implies that if *G* is 4-edge-connected, then for any  $X \subseteq E(G)$  with  $|X| \leq 4$ , *G* has a spanning Eulerian subgraph *H* such that  $X \subseteq E(H)$ . Here we have:

**Theorem 3.7.** Let G be 4-edge-connected graph. Let  $X \subseteq E(G)$  with  $|X| \leq 5$ . Let  $G_X$  be the graph obtained from G by subdividing each edge in X. Let  $D_i = \{v \in V(G'_X) | d(v) = i\}$   $(i \geq 2)$ . Then one of the following holds

- (i)  $G_X$  is collapsible, or
- (ii) X contains an edge cut  $X_1$  with  $|X_1| = t \ge 4$  such that  $G X_1$  has only two components ( $H_1$  and  $H_2$ ), which are collapsible. Furthermore,  $G_X$  is contractible to  $K_{2,t}$  by contracting  $H_1$  and  $H_2$  into the two degree t vertices in  $K_{2,t}$ , or
- (iii)  $G'_X$  is an Eulerian graph with  $V(G'_X) = D_2 \cup D_4$  and  $|D_2| = 5$ .

**Proof.** Let  $G'_X$  be the reduction of  $G_X$ . If  $G'_X = K_1$ , then  $G_X$  is collapsible and we are done for this case. In the following we will assume that  $G'_X$  is not trivial. Since G is 4-edge-connected,  $G_X$  is 2-edge-connected. Since  $\kappa(G'_X) \ge \kappa(G_X)$ ,  $G'_X$  is 2-edge-connected.

Case 1.  $F(G'_X) \leq 2$ .

By Theorem 2.5, and  $\kappa'(G_X) \ge 2$ ,  $G'_X = K_{2,t}$  for some  $t \ge 2$ . By Lemma 2.7,  $|X| \ge t \ge 4$ . Hence, (ii) of Theorem 3.7 holds.

Case 2.  $F(G'_{\mathbf{v}}) \ge 3$ . Since G is 4-edge-connected and  $|X| \leq 5$ , by (i) and (iii) of Lemma 2.7,  $V_r = D_2$  and

 $10 + |D_5| + \dots + \leq 2|V_r| \leq 2|X| \leq 10.$ 

This implies that  $|D_i| = 0$  for all  $i \ge 5$  and  $|D_2| = 5$ . Therefore, each vertex in  $V(G'_X)$  has degree 2 or 4. Hence,  $G'_X$  is Eulerian and  $|D_2| = 5$ .

**Corollary 3.8.** Let G be a 4-edge-connected graph. Let  $X \subseteq E(G)$  with  $|X| \leq 5$ . Let  $G_X$  be the graph obtained from *G* by subdividing each edge in *X*. Then either *G* has a spanning Eulerian subgraph *H* such that  $X \subseteq E(H)$ , or  $G_X$  is contractible to  $K_{2,5}$  in such a way that each path between the two vertices of degree 5 is obtained by subdividing an edge in X.

**Proof.** This follows from Theorem 3.7 and Lemma 2.9.  $\Box$ 

#### 4. The *r*-edge-Eulerian-connected graphs

We will need the following lemma.

**Lemma 4.0.** Let G be a 3-edge-connected graph. Let  $X \subseteq E(G)$  and let e',  $e'' \in E(G)$ . Let  $X_0 = X \cup \{e', e''\}$  and let  $G_{X_0}$  be the graph obtained from G by subdividing each edge in  $X_0$ . Suppose that  $G'_{X_0} = K_{2,t}$  where  $t \ge 3$ . If t > |X|, then G has a spanning (e', e'')-trail H such that  $X \subseteq E(H)$ .

**Proof.** Let u and v be the two vertices in  $K_{2,t}$  with d(u) = d(v) = t. By Lemma 2.7, there is an edge set  $X_1 \subseteq X_0$ such that each length 2 path between u and v in  $K_{2,t}$  is obtained by subdividing an edge in  $X_1$ . Then  $|X_1| = t$ . Let  $E_1 = E(G'_{X_0}) = E(K_{2,t})$ . By Lemma 2.7,  $G_{X_0} - E_1$  has two collapsible subgraphs ( $H_1$  and  $H_2$ ) such that  $V(G_{X_0}) = V(H_1) \cup V(H_2) \bigcup_{e \in X_1} \{v(e)\}$ . Let  $e' = x'_0 y'_0$ ,  $e'' = x''_0 y''_0$  and let  $x'_0, x''_0 \in V(H_1)$  and  $y'_0, y''_0 \in V(H_2)$ . Since t > |X|, at least one of the edges in  $\{e', e''\}$  is included in  $X_1$ . For each  $e \in \{e', e''\}$ ,  $P_e$  is defined as a path obtained by subdividing edge e.

For each  $H_i$ , (i = 1, 2), define

 $U_0(H_i) = \{v \in V(H_i) : v \text{ is incident with odd number of edges in } E_1 - \{P_{e'}, P_{e''}\}\}$ 

Note that  $|U_0(H_1)|$  is odd if and only if  $|U_0(H_2)|$  is odd. Since  $H_i$  is collapsible, for any even subset  $R_i \subseteq V(H_i)$ , there is a spanning connected subgraph  $\Gamma_i$  with  $O(\Gamma_i) = R_i$  (i = 1, 2). In the following we will show that a spanning (v(e'), v(e'))-trail  $\Gamma$  can be constructed from  $\Gamma_1$  and  $\Gamma_2$  by adding all the edges in  $E_1$  and an edge  $e_{\Gamma_1}$  to connect v(e')(or an edge  $e_{\Gamma_2}$  to connect v(e''), or both) such that  $O(\Gamma) = \{v(e'), v(e'')\}$ .

*Case* 1. Both e' and e'' are in  $X_1$ .

Note that G may not be simple and we may have three possible situations:

- (a)  $x'_0 = x''_0$  and  $y'_0 = y''_0$ ,
- (b)  $x'_0 = x''_0$  and  $y'_0 \neq y''_0$ , (c)  $x'_0 \neq x''_0$  and  $y'_0 \neq y''_0$ .

The following Tables 1–3 show the selections of the even subset  $R_i \subseteq V(H_i)$  for  $\Gamma_i$  and  $e_{\Gamma_i}$  (i = 1, 2) for all possible cases.

For each case with the selection of  $R_1$ ,  $R_2$ ,  $e_{\Gamma_1}$  and  $e_{\Gamma_2}$ , define

$$\Gamma = G_{X_0}[E(\Gamma_1) \cup E(\Gamma_2) \cup E_1 \cup \{e_{\Gamma_1}, e_{\Gamma_2}\}].$$

By the definition of  $\Gamma$ ,  $V(\Gamma) = V(\Gamma_1) \cup V(\Gamma_2) \bigcup_{e \in X_1} \{v(e)\} \cup \{v(e'), v(e'')\}$ , and v(e') and v(e'') have degree 1 in  $\Gamma$ . Since  $\Gamma_i$  is a connected spanning subgraph of  $H_i$ ,  $V(\Gamma_i) = V(H_i)$  (i = 1, 2).  $\Gamma_1$  and  $\Gamma_2$  are connected by the paths in  $E_1$ , and v(e') and v(e'') are connected to  $\Gamma_i$  by  $e_{\Gamma_i}$ . Thus,  $V(\Gamma) = V(G_{X_0})$  and  $\Gamma$  is a connected spanning subgraph

Table 1 When  $x'_0 = x''_0$  and  $y'_0 = y''_0$ , let  $x_0 = x'_0 = x''_0$  and  $y_0 = y'_0 = y''_0$ 

$ U_0(H_1) $	$x_0$ and $y_0$	<i>R</i> <sub>1</sub>	<i>R</i> <sub>2</sub>	$e_{\Gamma_1}$	$e_{\Gamma_2}$
Odd	$\begin{aligned} x_0 &\in U_0(H_1),  y_0 \in U_0(H_2) \\ x_0 &\notin U_0(H_1),  y_0 \in U_0(H_2) \\ x_0 &\in U_0(H_1), y_0 \notin U_0(H_2) \\ x_0 &\notin U_0(H_1),  y_0 \notin U_0(H_2) \end{aligned}$	$U_{0}(H_{1}) - x_{0}$ $U_{0}(H_{1}) \cup \{x_{0}\}$ $U_{0}(H_{1}) - x_{0}$ $U_{0}(H_{1}) \cup \{x_{0}\}$	$U_{0}(H_{2}) - y_{0}$ $U_{0}(H_{2}) - y_{0}$ $U_{0}(H_{2}) \cup \{y_{0}\}$ $U_{0}(H_{2}) \cup \{y_{0}\}$	$     x_0 v(e')      x_0 v(e')      x_0 v(e')      x_0 v(e') $	$v(e'')y_0  v(e'')y_0  v(e'')y_0  v(e'')y_0$
Even		$U_0(H_1)$	$U_0(H_2)$	$x_0v(e')$	$x_0v(e'')$

Table 2		
When $x'_0 = x''_0$ and $y'_0$	$\neq y_0''$ , let $x_0$	$=x_0'=x_0''$

$ U_0(H_1) $	$x_0$ , and $y_0''$	<i>R</i> <sub>1</sub>	<i>R</i> <sub>2</sub>	$e_{\Gamma_1}$	$e_{\Gamma_2}$
Odd	$\begin{aligned} x_0 &\in U_0(H_1), y_0'' \in U_0(H_2) \\ x_0 &\in U_0(H_1), y_0'' \notin U_0(H_2) \\ x_0 &\notin U_0(H_1), y_0'' \in U_0(H_2) \\ x_0 &\notin U_0(H_1), y_0'' \notin U_0(H_2) \end{aligned}$	$U_{0}(H_{1}) - x_{0}$ $U_{0}(H_{1}) - x_{0}$ $U_{0}(H_{1}) \cup \{x_{0}\}$ $U_{0}(H_{1}) \cup \{x_{0}\}$	$\begin{array}{l} U_{0}(H_{2})-y_{0}''\\ U_{0}(H_{2})\cup\{y_{0}''\}\\ U_{0}(H_{2})-y_{0}''\\ U_{0}(H_{2})-y_{0}''\\ U_{0}(H_{2})\cup\{y_{0}''\} \end{array}$	$     x_0 v(e')      x_0 v(e')      x_0 v(e')      x_0 v(e') $	$v(e'')y_0'' \\ v(e'')y_0'' \\ v(e'')y_0'' \\ v(e'')y_0'' \\ v(e'')y_0''$
Even		$U_0(H_1)$	$U_0(H_2)$	$x_0v(e')$	$x_0v(e'')$

Table 3 When  $x'_0 \neq x''_0$  and  $y'_0 \neq y''_0$ 

$ U_0(H_1) $	$x'_0$ , and $y''_0$	<i>R</i> <sub>1</sub>	<i>R</i> <sub>2</sub>	$e_{\Gamma_1}$	$e_{\Gamma_2}$
Odd	$\begin{array}{l} x_0' \in U_0(H_1),  y_0'' \in U_0(H_2) \\ x_0' \in U_0(H_1),  y_0'' \notin U_0(H_2) \\ x_0' \notin U_0(H_1),  y_0'' \in U_0(H_2) \\ x_0' \notin U_0(H_1),  y_0'' \notin U_0(H_2) \end{array}$	$U_{0}(H_{1}) - x'_{0}$ $U_{0}(H_{1}) - x'_{0}$ $U_{0}(H_{1}) \cup \{x'_{0}\}$ $U_{0}(H_{1}) \cup \{x'_{0}\}$	$U_{0}(H_{2}) - y_{0}''$ $U_{0}(H_{2}) \cup \{y_{0}''\}$ $U_{0}(H_{2}) - y_{0}''$ $U_{0}(H_{2}) \cup \{y_{0}''\}$	$     x'_{0}v(e')      x'_{0}v(e')      x'_{0}v(e')      x'_{0}v(e') $	$v(e'')y_0''$ $v(e'')y_0''$ $v(e'')y_0''$ $v(e'')y_0''$
Even	$ \begin{array}{l} x_0' \in U_0(H_1), x_0'' \in U_0(H_1) \\ x_0' \notin U_0(H_1), x_0'' \in U_0(H_1) \\ x_0' \in U_0(H_1), x_0'' \notin U_0(H_1) \\ x_0' \notin U_0(H_1), x_0'' \notin U_0(H_1) \end{array} $	$U_{o}(H_{1}) - \{x'_{0}, x''_{0}\} \\ (U_{o}(H_{1}) - \{x''_{0}\}) \cup \{x'_{0}\} \\ (U_{o}(H_{1}) - \{x'_{0}\}) \cup \{x''_{0}\} \\ U_{o}(H_{1}) \cup \{x'_{0}, x''_{0}\}$	$U_{0}(H_{2}) \\ U_{0}(H_{2}) \\ U_{0}(H_{2}) \\ U_{0}(H_{2})$	$\begin{array}{c} x_{0}^{\prime}v(e^{\prime}) \\ x_{0}^{\prime}v(e^{\prime}) \\ x_{0}^{\prime}v(e^{\prime}) \\ x_{0}^{\prime}v(e^{\prime}) \end{array}$	$\begin{array}{c} x_0''v(e'') \\ x_0''v(e'') \\ x_0''v(e'') \\ x_0''v(e'') \end{array}$

of  $G_{X_0}$ . To show that  $O(\Gamma) = \{v(e'), v(e'')\}$ , we can check each case listed in Tables 1–3. For instance, with the cases in Table 1, if  $v \notin R_1 \cup R_2$ , v has even degree in  $\Gamma_1$  or  $\Gamma_2$  or v has degree 2 as a vertex obtained by subdividing an edge in  $X_1$ . If  $v \notin R_1$  and  $v \neq x_0$  (or  $v \notin R_2$  and  $v \neq y_0$ ), then since odd number of edges incident with v in  $E_1$  are added, v has an even degree in  $\Gamma$ . If  $v = x_0$  (or  $y_0$ ), by the definition of  $e_{\Gamma_1}$  and  $e_{\Gamma_2}$ ,  $x_0$  has an even degree in  $\Gamma$ . Hence,  $O(\Gamma) = \{v(e'), v(e'')\}$ , and  $\Gamma$  is a spanning (v(e'), v(e''))-trail in  $G_{X_0}$ . By Lemma 2.9, G has a spanning (e', e'')-trail containing X.

*Case* 2. One of e' and e'' is in  $X_1$  (say  $e' \in X_1$ ).

Since  $e'' \notin X_1$ , we may assume that the path obtained by subdividing e'' is in  $H_1$ . Then  $v(e'') \in V(H_1)$ . For this case, we only need to choose  $e_{\Gamma_1}$  to connect v(e') in  $\Gamma$ .

For each case in Table 4, define

 $\Gamma = G_{X_0}[E(\Gamma_1) \cup E(\Gamma_2) \cup E_1 \cup \{e_{\Gamma_1}\}].$ 

Therefore,  $\Gamma$  is a spanning connected subgraph of  $G_{X_0}$  such that  $O(\Gamma) = \{v(e'), v(e'')\}$ . The Lemma is proved.  $\Box$ 

In [14], Zhan proved the following:

**Theorem 4.1** (*Zhan* [14]). If G is a 4-edge-connected graph, then for any edges  $e_1, e_2 \in E(G)$  there is a spanning  $(e_1, e_2)$ -trail in G.

$e' \in X_1$ , and $v(e'') \in V(H_1)$				
$ U_0(H_1) $	$x'_0$ , and $y'_0$	$R_1$	<i>R</i> <sub>2</sub>	$e_{\Gamma_1}$
Odd	$y'_0 \in U_0(H_2)$ $y'_0 \notin U_0(H_2)$	$U_{o}(H_{1}) \cup \{v(e'')\} \\ U_{o}(H_{1}) \cup \{v(e'')\}$	$U_{0}(H_{2}) - y'_{0}$ $U_{0}(H_{2}) \cup \{y'_{0}\}$	$v(e')y'_0 \\ v(e')y'_0$
Even	$\begin{array}{l} x_0' \in U_0(H_1) \\ x_0' \notin U_0(H_1) \end{array}$	$ \begin{array}{l} (U_0(H_1) - \{x'_0\}) \cup \{v(e'')\} \\ U_0(H_1) \cup \{x'_0, v(e'')\} \end{array} $	$U_{0}(H_{2})$ $U_{0}(H_{2})$	$\begin{array}{c} x_0' v(e') \\ x_0' v(e') \end{array}$

Theorem 4.1 can be improved.

Table 4

**Theorem 4.2.** Let  $r \in \{3, 4\}$ . If G is an (r + 1)-edge-connected graph, then for any  $X \subseteq E(G)$  with  $|X| \leq r - 1$ , and for any  $e_1, e_2 \in E(G)$ , G has a spanning  $(e_1, e_2)$ -trail H in G such that  $X \subseteq E(H)$ .

**Proof.** Let  $X_0 = X \cup \{e_1, e_2\}$ . Let  $G_{X_0}$  be the graph obtained from *G* by subdividing each edge in  $X_0$ . Since  $r \in \{3, 4\}$ ,  $k = \lfloor (r+1)/2 \rfloor = 2$ . Then  $|X_0| \leq |X| + 2 \leq r+1 = (r+1)+k-2$ . By Theorem 3.1', either  $G_{X_0}$  is collapsible or  $G_{X_0}$  is contractible to  $K_{2,t}$  with  $t \geq r$ . If  $G_{X_0}$  is collapsible, then by Lemma 2.9, *G* has a spanning  $(e_1, e_2)$ -trail containing *X*. If  $G_{X_0}$  is contractible to  $K_{2,t}$  with  $t \geq 4$ , since  $t \geq r > |X|$ , by Lemma 4.0, *G* has a spanning  $(e_1, e_2)$ -trail containing the edge set *X*.  $\Box$ 

For graphs with edge-connectivity at least 5, we have

**Theorem 4.3.** Let G be an (r + 1)-edge-connected graph  $(r \ge 4)$ . Let  $X \subseteq E(G)$  with  $|X| \le r$ . Then G is an r-edge-Eulerian-connected.

**Proof.** Let  $e_1$  and  $e_2$  be two arbitrary edges in G and let  $X_0 = X \cup \{e_1, e_2\}$ . Let  $G_{X_0}$  be the graph obtained from G by subdividing each edge in  $X_0$ .

Case 1.  $r \ge 5$ .

Then  $r + 1 \ge 6$ , and so  $k = \lfloor (r + 1)/2 \rfloor \ge 3$ . Then  $|X_0| \le |X| + 2 \le r + 2 \le (r + 1) + k - 2$ . By Theorem 3.1', either  $G_{X_0}$  is collapsible or  $G_{X_0}$  is contractible to  $K_{2,t}$  with  $|X_0| \ge t \ge (r + 1)$ . By Lemma 2.9 and Lemma 4.0, both cases imply that *G* has a spanning  $(e_1, e_2)$ -trail *H* such that  $X \subseteq E(H)$ . Theorem 4.3 is proved for this case.

Case 2. 
$$r = 4$$

Then *G* is 5-edge-connected and  $|X_0| \leq 6$ . Let  $G'_{X_0}$  be the reduction of  $G_{X_0}$ . If  $F(G'_{X_0}) \leq 2$ , then  $G_{X_0}$  is either collapsible or contractible to  $K_{2,t}$  with  $t \geq (r+1)$  and so we are done. Next we assume that  $F(G'_{X_0}) \geq 3$ .

**Claim.** If  $v \in D_2 \subseteq V(G'_{X_0})$ , then the degree of each of the two neighbors of v is greater than 2.

Since  $\delta(G) \ge \kappa'(G) \ge 5$ , each vertex of degree 2 in  $G'_{X_0}$  is obtained by subdividing an edge in  $X_0$ . If a degree vertex has a neighbor which is also degree, then this will contradict to the definition of  $G_{X_0}$ .

By Lemma 2.8, we have

$$|V(G'_{X_0})| + 10 \leqslant 3|D_2| \leqslant 3|X_0|. \tag{8}$$

If  $|D_2| \leq 5$ , then by (8),  $|V(G'_{X_0})| \leq |D_2| \leq 5$ , contrary to the claim above. Therefore,  $|D_2| = |X_0| = 6$ . By (8) and  $|D_2| = 6$ ,

 $|V(G'_{X_0})| \leq 8.$ 

Therefore,  $G'_{X_0}$  is a 2-edge-connected graph with 6 vertices of degree 2 and at most two vertices of degree at least 5. By the claim above, vertices of degree 2 are not adjacent to each other. Therefore,  $G'_{X_0} = K_{2,6}$ , contrary to  $F(G'_{X_0}) \ge 3$ . The theorem is proved.  $\Box$ 

Let *r* be an integer. Theorem 4.2 shows that if *G* is 4-edge-connected, then *G* is 2-edge-Eulerian-connected. If  $r \ge 4$  and if *G* is (r + 1)-edge-connected, then *G* is *r*-edge-Eulerian-connected. Combining Theorems 4.2, 4.3 and 3.0, we have:

**Corollary 4.4.** Let  $r \ge 0$  be an integer. Then

$$\xi(r) = \begin{cases} 4, & 0 \le r \le 2, \\ r+1, & r \ge 4. \end{cases}$$

**Remark.** The case  $\xi(3)$  is still open. Theorem 4.2 implies that if G is 5-edge-connected, then G is 3-edge-Eulerianconnected, and so  $\xi(3) \leq 5$ . We conjecture that  $\xi(3) = 4$ . The following theorem provides some supports for this conjecture.

**Theorem 4.5.** Let G be a 4-edge-connected graph and let  $X \subseteq E(G)$  with  $|X| \leq 3$ . For any two adjacent edges e' and e'', G has a spanning (e', e'')-trail H such that  $X \subseteq E(H)$ .

**Proof.** Let  $X_0 = X \cup \{e', e''\}$ . Let  $G_{X_0}$  be the graph obtained from G by subdividing each edge in  $X_0$ . Let v(e') and v(e'') be the two vertices obtained in the process of subdividing e' and e''. If  $G_{X_0}$  is collapsible, then  $G_{X_0}$  has a spanning connected subgraph H such that  $O(H) = \{v(e'), v(e'')\}$ . By Lemma 2.9, G has a spanning (e', e'')-trail containing X. We are done in this case. Next, we assume that  $G_{X_0}$  is not collapsible.

Let  $G'_{X_0}$  be the reduction of  $G_{X_0}$ . By Theorem 3.7, either  $G'_{X_0} = K_{2,t}$  with  $t \ge 4$  or  $G'_{X_0}$  is Eulerian with  $V(G_{X_0}) = D_2 \cup D_4$  and  $|D_2| = 5$ , where  $D_i$  is the set of vertices of degree *i* in  $G'_{X_0}$ . If  $G'_{X_0} = K_{2,t}$  with  $t \ge 4$ , then by Lemma 4.0, *G* has a spanning (e', e'')-trail *H* such that  $X \subseteq E(H)$ . We are done for this case.

For the case that  $G'_{X_0}$  is Eulerian, let v be the vertex incident with both e' and e''. Let  $e_1 = v(e')v$  and  $e_2 = v(e'')v$ . Then  $G'_{X_0} - \{e_1, e_2\}$  is connected. Otherwise,  $\{e', e''\}$  is an edge cut of G, contrary to that G is 4-edge-connected. Therefore,  $G'_{X_0} - \{e_1, e_2\}$  is a connected graph with only two odd degree vertices at v(e') and v(e''). Let  $U_4 = \{u \in D_4 : u \text{ is a non-trivial contraction}\}$ . For each vertex  $u \in U_4$ , let H(u) be the preimage of u in  $G_{X_0}$ . Then H(u) is collapsible. Let

 $V_u = \{x \in V(H(u)) : x \text{ is incident with odd number of edges in } G'_{X_0} - \{e_1, e_2\}\}.$ 

Since d(u) in  $G'_{X_0} - \{e_1, e_2\}$  is even,  $|V_u|$  is even or 0. Since H(u) is collapsible, H(u) has a spanning connected subgraph  $\Gamma_u$  such that  $O(\Gamma_u) = V_u$ . Let  $E_0 = E(G_{X_0}) - \{e_1, e_2\}$  and let

$$\Gamma = G_{X_0} \left[ \bigcup_{u \in U_4} E(\Gamma_u) \cup E_0 \right].$$

Then  $\Gamma$  is a spanning connected subgraph of  $G_{X_0}$  such that  $O(\Gamma) = \{v(e'), v(e'')\}$ . Therefore,  $G_{X_0}$  has a spanning (v(e'), v(e''))-trail. By Lemma 2.9, G has a spanning (e', e'')-trail containing X. The proof is complete.  $\Box$ 

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