Air Force Institute of Technology AFIT Scholar

Theses and Dissertations

Student Graduate Works

3-1994

A Genetic Algorithm Approach to Automating Satellite Range Scheduling

Donald A. Parish

Follow this and additional works at: https://scholar.afit.edu/etd

Part of the Space Vehicles Commons

Recommended Citation

Parish, Donald A., "A Genetic Algorithm Approach to Automating Satellite Range Scheduling" (1994). *Theses and Dissertations*. 6770. https://scholar.afit.edu/etd/6770

This Thesis is brought to you for free and open access by the Student Graduate Works at AFIT Scholar. It has been accepted for inclusion in Theses and Dissertations by an authorized administrator of AFIT Scholar. For more information, please contact AFIT.ENWL.Repository@us.af.mil.

AFIT/GOR/ENS/94M-10



A GENETIC ALGORITHM APPROACH TO

a . 1 🌤

AUTOMATING SATELLITE RANGE SCHEDULING

THESIS Donald Arthur Parish Captain, USAF

AFIT/GOR/ENS/94M-10



DITC QUALLEY INCRED 3

Approved for public release; distribution unlimited

4 21 050

94



The views expressed in this thesis are those of the author and do not reflect the official policy or position of the Department of Defense or the U.S. Government.

•*

Accesion	For			
NTIS CRA&I DTIC TAB Unannounced Justification				
1	By Distribution /			
A	vailabilit	y Codes		
Dist	Dist Avail and or Special			
A-1				

THESIS APPROVAL

STUDENT: Donald A. Parish, Capt, USAF

THESIS TITLE: A GENETIC ALGORITHM APPROACH TO AUTOMATING SATELLITE RANGE SCHEDULING

DEFENSE DATE: 1 March 94

COMMITTEE:

Name/Title/Department

Advisor: James W. Chrissis, PhD, P.E. Associate Professor of Operations Research **Department of Operational Sciences** School of Engineering

Reader: Gary B. Lamont, PhD **Professor of Electrical Engineering** Department of Electrical and Computer Engineering School of Engineering

James T. Moore, Lt Col, USAF, PhD Assistant Professor of Operations Research **Department of Operational Sciences** School of Engineering

Signature

amer 1. Moon

CLASS: GOR-94M

Reader:

AFIT/GOR/ENS/94M-10

A GENETIC ALGORITHM APPROACH TO AUTOMATING SATELLITE RANGE SCHEDULING

THESIS

Presented to the Faculty of the Graduate School of Engineering of the Air Force Institute of Technology Air University In Partial Fulfillment of the Requirements for the Degree of Master of Science in Operations Research

> Donald Arthur Parish, B.S. Captain, USAF

> > March, 1994

Approved for public release; distribution unlimited

Acknowledgements

I would like to thank everyone who helped with this thesis. Of special note were my advisor Dr. James W. Chrissis for his solid recommendations, and my readers:

Dr. Gary B. Lamont and Lt Col James T. Moore for their patience and understanding. I would also like to thank those whose work this research was based on: Capt Timothy D. Gooley, who was also my sponsor at AFIT, and Capt Stanley M. Schalck, who answered many of my questions about satellite range scheduling and data processing. Thanks are also due to my classmates and instructors who made the journey here an interesting one. Not least, I am grateful for the support and encouragement of my wife Lisa. I owe her much attention and a trip back to Nebraska.

Donald Arthur Parish

Table of Contents

.•

	Page
Acknowledgements	ii
List of Figures	vii
List of Tables	viii
Abstract	ix
I. Introduction	1
Problem Description	2
Previous Solution Efforts	3
Computational Complexity	4
Research Objective	4
Overview	5
II. Literature Review	6
Genetic Algorithms	6
Introduction	6
Simple Genetic Algorithm	6
Why GAs Work	9
Permutation Genetic Algorithms	10
Traveling Salesman Problem	11
Blind Traveling Salesman Problem	11
Order-based Crossover Operators	12
-	
Ordering Schema	13
Genetic Algorithms in Scheduling	15
Direct Chromosome Representation	15

		Page
Indirect Representation	 ••••	16
Scheduling Examples	 	16
Other Solution Efforts for Satellite Range Scheduling	 • • • •	18
Arbabi's Approach	 • • • •	18
Gooley's Approach	 • • • •	18.
Schalck's Approach	 • • • •	19
Summary	 • • • •	19
III. Solution Methodology for Satellite Range Scheduling	 • • • •	21
Formulation	 	2 1
Definition of the Problem	 • • • •	21
Small Problem	 ••••	2 1
Mixed-Integer Programming Approach	 • • • •	23
Scheduling as a Sequencing Problem	 • • • •	23
Hybrid Approach	 • • • •	25
Algorithm Design and Implementation	 • • • •	25
Overview	 • • • •	25
Assumptions	 • • • •	26
Data Processing	 • • • •	26
Data Structures	 	27
Schedule Builder Program	 	28
Genetic Algorithm Implementation	 • • • •	29
GENITOR Parameter Settings	 • • • •	30
Day One Results	 • • • •	31
Overall Implementation	 • • • •	33
Summary	 	33

	Page
IV. Results	3 5
Test Set	35
Experimental Procedure	36
Week One Results	37
Breakdown by Support Type	38
Additional Runs	38
Support Priority	39
Two-Day Scheduling	40
Summary	41
V. Conclusions and Recommendations	44
Conclusions	44
Recommendations	45
Minor Improvements	45
Extensions	45
Summary	46
Appendix A. Modified GENITOR Program Code	47
Genetic Algorithm: GENITOR	47
Modified Main program Code	48
Ga global.h Include files	51
Position Crossover Operator	53
Appendix B. Evaluation Function Code	57
Appendix C. Schedule Builder	63
Appendix D. Satellite Range Scheduling Data Processing	69
ASTRO Data	69
Requests and Visibility Windows	69

.

	Page
LREQ.PAS	70
HREQ.PAS	70
TOL.PAS	70
CROSS2.PAS	70
RTS.PAS	70
Prepare Time Window Data for GENITOR	70
ppendix E. Schedules for Week One Data	73
Day 1	74
Day 2	75
Day 3	77
Day 4	79
Day 5	81
Day 6	83
Day 7	84
ibliography	87
ita	89

. -

List of Figures

....

、 - [•]

Figure		Page
1.	Genetic Algorithm Procedure	7
2.	Example of Crossover	8
3.	Example of Mutation	8
4.	Sample TSP tours	11
5.	Small Schedule	22
6.	Schedule Builder Flowchart	25
7.	Schedule Builder program flow	29
8.	Day 1 Data Convergence	32
9.	Vary Population Size for Day 1 Data	33
10.	Vary Selective Pressure for Day 1 Data	34
11.	Random Sequence Generation for Day 1 Data	34
12 .	Week One Data Results	42
13.	Low-Altitude Supports given Priority for Day 1 Data	43

List of Tables

.

Table		Page
1.	Small Schedule Time Windows	22
2 .	Small Problem Schedule	24
3.	Small Schedule Time Windows	27
4.	Results for Week One Data	37
5.	Low-altitude Results for Week One Data	38
6.	High-altitude Results for Week One Data	39
7.	Week One Data with Low-Altitude Supports given Priority	39
8.	Low-altitude Results with Low-Altitude Supports given Priority	40
9.	High-Altitude Results with Low-Altitude Supports given Priority	40

Abstract

Satellite range scheduling involves scheduling satellite supports in which a satellite and a specific remote tracking station communicate with each other within a specified time window. As the number of satellite supports continue to increase, more pressure is placed on the current manual system to generate schedules efficiently. Previous research efforts focused on heuristic and mixed-integer programming approaches which may not produce the best results. The objective of this research was to determine if a genetic algorithm approach to automating the generation of 24 hour schedules was competitive with other methods. The goal was to schedule as many supports as possible without conflict.

The genetic algorithm approach attempted to find the best priority ordering of support requests, and then used a schedule builder program to build schedules based on simple rules. A schedule was produced for seven days of representative satellite range data with slightly better results compared to earlier results using a mixed-integer programming formulation. Based on the reported results, the genetic algorithm approach presented in this research appears to be a competitive approach for generating 24-hour satellite range schedules.

A GENETIC ALGORITHM APPROACH TO AUTOMATING SATELLITE RANGE SCHEDULING

I. Introduction

Scheduling is the allocation of resources over time to perform a collection of tasks (2:2). Many important scheduling problems exist that are of interest to the Air Force. Examples include allocation of test range resources, airlift scheduling problems, and satellite communication support scheduling. Many standard scheduling problems may be formulated as mixed-integer programming problems where some decision variables take on integer values while other variables, such as time, have continuous values. While a mathematical programming formulation generally guarantees an optimal, or best possible, solution, the computation times required to find exact optimal solutions may be prohibitive for practical-sized problems. This is because many scheduling problems belong to the class of NP-complete problems (6:4).

In such NP-complete problems, the solution time may increase exponentially with the number of variables. For example, a problem with 10 times the number of variables as an original problem might take an order of 2^{10} times longer to solve. Because of these possible long solution times, heuristic techniques are often used to find good solutions in a reasonable amount of time. Heuristics are procedures which do not guarantee optimal solutions. They usually, but not always, provide feasible solutions, often require human ingenuity in their development, and are often quite problem-specific.

Heuristic methods for solving scheduling problems that are less dependent on the specific problem could be useful in solving problems of Air Force interest. This would be especially true when the constraints of the problem are difficult to put in mathematical form. One such method for finding good solutions to scheduling and other optimization problems is called a *genetic algorithm*.

Genetic algorithms (GAs) are artificial intelligence search methods based on the idea of natural selection and evolution. Initially developed by John Holland at the University of Michigan (10), applications include problems in optimization and machine learning. Although initially applied to function optimization problems, genetic algorithms have also been applied to scheduling and combinatorial optimization problems. The main strength of genetic algorithms is their ability to quickly explore a large space of possible solutions for good, if not optimal, solutions. They differ from many traditional algorithms as their search is based only on the overall evaluation of a set of parameters. As such, they do not need to rely on derivative information to proceed. Genetic algorithms are also able to find solutions where multiple optimal solutions exist (7:2-5).

Problem Description

One particular scheduling problem of interest to the Air Force is the Satellite Range Scheduling (SRS) problem. The Air Force Satellite Control Network (AFSCN) must schedule over 300 command and control communications per day between nine remote tracking stations (RTSs) and approximately 100 satellites (8:1-3). Each communication between a satellite and a RTS is called a *support*. The satellite supports are necessary to maintain command and control, tracking, and system tests of the satellites (11:10). The scheduling of these supports must take place in a *time window* because each RTS is geographically separated from the others, and can only "see" each satellites than for medium- or highaltitude satellites, as the low-altitude satellites pass out of the line-of-sight of the RTSs much more quickly than higher-altitude satellites. The longer time window of the higheraltitude satellites makes scheduling them less difficult than scheduling the low-altitude satellites.

Each Mission Control Complex (MCC) is responsible for the health, status, and orbital control of a subset of the total satellites. The composition of the subset depends on the mission type of the satellite. A MCC determines the length and required time windows for each support request for their satellites, as well as which RTSs can serve each request. These time windows may be more restricted than the physical visibility limits due to mission scheduling requirements. For example, a satellite may need a communication every hour after an operation is performed. All requests from the individual MCCs are passed on to the range schedulers who are responsible for making an overall schedule. This schedule must ensure that each RTS antenna supports only one satellite at a time. The schedulers must also allow for a required turn-around time between supports to allow the RTS antennas to be reoriented. Downtime for RTS maintenance is also necessary. The resulting schedule is called the *initial* 24-hour schedule. The development of this schedule using genetic algorithms is the focus of this research.

The goal of the initial 24-hour scheduling process is to schedule as many communication supports as possible in a 24-hour period while satisfying the constraints of time windows, turnaround time, and scheduled maintenance down-time for the RTSs. If a conflict cannot be resolved by the range schedulers, they must de-conflict it with the MCCs and RTSs involved. If a support cannot be scheduled initially, it may be possible to facilitate its scheduling by changing the support requirements, by decreasing the turn-around time, or by altering other constraints. In the worst case, a support may not be scheduled. A good scheduling process minimizes the need for this deconfliction process. Although the current manual scheduling process schedules approximately 95-98% of the requested supports (8:5-2), the process is time-consuming, and could be streamlined by automating the development of the 24-hour schedule.

Previous Solution Efforts

Two recent thesis efforts have investigated automation of the 24-hour scheduling process. Gooley used both a mixed-integer programming (MIP) approach and heuristic scheduling methods to schedule satellite supports (8). As a follow-on research effort, Schalck improved Gooley's solution by reducing the number of integer variables necessary to define the problem (14). In both cases, the problem was too large to be solved in its entirety and had to be decomposed into smaller problems to find a solution in reasonable time.

Computational Complexity

Computational problems are often classified based on their complexity. According to Garey and Johnson (6:4), problems that are classified as NP-complete can take an inordinate amount of time to solve if the size of the problem is large enough. The general resource-constrained scheduling problem is in this class. According to Gooley, the SRS problem is a type of resource-constrained scheduling problem (8:2-9), and a fast (polynomial) solution cannot be guaranteed. Genetic algorithms, which have had some success in finding good solutions to other NP-complete scheduling problems, may be useful when applied to the SRS problem.

The mixed-integer programming approach can find optimal solutions for the satellite range scheduling problem when the number of support requests is small. However, when a larger problem is decomposed into smaller subproblems, a global optimal solution may no longer be guaranteed. A genetic algorithm approach may be expected to produce schedules that are at least as good while meeting the requirement of a short solution time since it would attempt to generate a solution to the entire problem. Also, a genetic algorithm approach can be more flexible in handling special case scheduling requests. These issues are addressed in this research.

Research Objective

The overall research objective is to determine whether a genetic algorithm-based solution methodology can effectively be applied to the satellite range scheduling problem. Successful solutions must:

- 1. Generate feasible 24-hour schedules which schedule the greatest number of support requests possible.
- 2. Find solutions quickly. A short solution time is useful in rescheduling as requirements change.

A secondary objective is to explore the scheduling of special-case requirements. These include supports which require simultaneous support on multiple RTSs and those that must be scheduled at fixed intervals. Rescheduling of satellite supports is also a consideration.

Overview

Chapter II is a summary of current literature relevant to genetic algorithm applications to scheduling problems and includes a review of previous work on the satellite range scheduling problem. Chapter III explains the methodology used in constructing a genetic algorithm solution to the satellite range scheduling process and details the implementation of the solution. Chapter IV presents the results and compares them to past efforts. Finally, conclusions and recommendations for future work are presented in Chapter V.

II. Literature Review

This chapter summarizes current literature on genetic algorithms (GAs) as related to scheduling problems. It begins with the development of genetic algorithms. Next, the extension of genetic algorithms to combinatorial optimization problems is reviewed. Discussion then turns to applications of genetic algorithms in scheduling problems, especially to problems with similarities to the satellite range scheduling problem. Finally, past research efforts in satellite range scheduling are discussed.

Genetic Algorithms

Introduction. Genetic algorithms were developed by John Holland as part of his artificial intelligence research on how artificial adaptive systems can evolve, or change, in response to their environment in order to solve problems (10). His ideas were based on the biological theory of evolution where variations in *chromosomes*, or genetic codes, result in different traits of the individual. These traits of an individual in turn result in a level of performance, or *fitness*, of an individual. Usually, the more fit individuals in a *population* survive from one generation to another and reproduce. Sexual reproduction by two individuals produces individuals with new chromosomes formed by a combination of the chromosomes of its parents. If a child inherits good parts of chromosomes from each of its parents, it has a higher level of performance than the parents. Through "survival of the fittest" and reproduction of high-fitness individuals in each generation, the population as a whole tends to evolve towards higher levels of fitness. The artificial version of this process is called a genetic algorithm. It does not seek to exactly simulate biological evolution, but the concepts of biological evolution are used to find good solutions to difficult problems.

Simple Genetic Algorithm. Most genetic algorithm work is based on the simple genetic algorithm as developed by Holland (10) and described by Michalewicz (13). Figure 1 illustrates the process and a brief review follows.

Coding. The first, and generally most difficult, step of any genetic algorithm is to choose a proper coding to map the problem solution space into a genetic string,

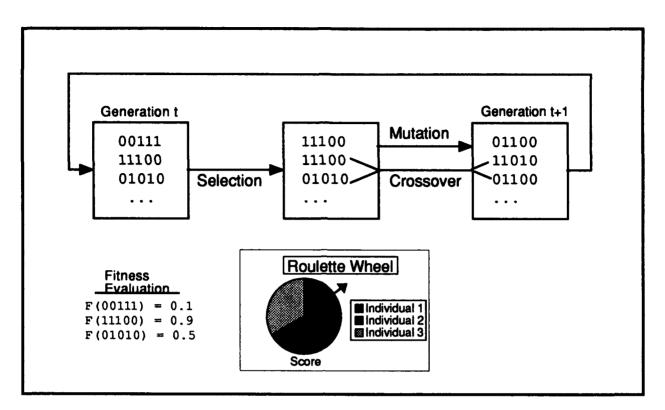


Figure 1. Genetic Algorithm Procedure

or chromosome, and to randomly create an initial population of individuals with varying strings. In the simple genetic algorithm, this coding is a binary string of zeros or ones. For function optimization, groups of binary digits are mapped so as to translate to a real number parameter representing the function value (13:19-20).

Evaluation and Selection. The strings in the population can be evaluated for their fitness relative to other strings in the population by entering their parameters into an evaluation or fitness function. The best strings reproduce by mating with each other to produce offspring for the next generation of the population. In the simple genetic algorithm, selection is governed by a "roulette wheel" selection operator. Each string has a probability of reproducing in proportion to the ratio of its fitness and the total fitness of the population. As in Figure 1, these ratios can be shown as pieces of a circular pie. New strings are selected for reproduction by randomly "spinning" the wheel. The strings whose proportion of the pie are greatest should, on average, be selected more often than the less-fit strings. Other selection rules can be used (13:62). These include selection by ranking, where instead of a continuous scale based on fitness the strings are ordered, or ranked, by their fitness. The higher-ranked strings are given a greater chance of reproducing than average or low-ranked strings. In this way, the more fit individuals tend to reproduce.

Crossover Operator. During mating, two strings swap part of their "genetic" material. The children of these matings have parts of each of their parent's string. The swapping of genetic material, called *crossover*, allows for new child strings to be created, combining good aspects of their parents. The resulting strings with above-average fitness tend to survive and prosper, while those with below-average fitness tend to die out. An example of crossover is shown in Figure 2.

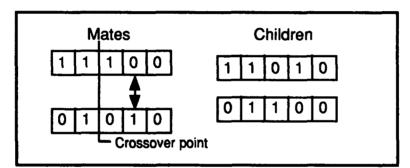


Figure 2. Example of Crossover

Mutation. Mutation randomly changes part of the string of a child to help maintain a diverse population. It is not as important as crossover in some applications since it merely acts as a type of random search (7:14).

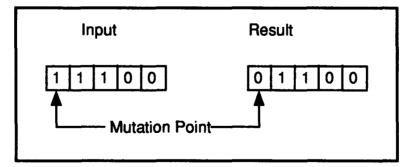


Figure 3. Example of Mutation

The general genetic algorithm thus consists of three steps which are repeated for each generation: evaluation/selection, crossover and mutation. These steps continue until termination conditions, such as some predetermined number of generations, are met. Often the genetic algorithm ends when the population converges: all strings evaluate to the same fitness (13:56). This string or strings may be the best answer, but a genetic algorithm may often converge to a suboptimum, in which case the population has prematurely converged. The major variables controlled by the experimenter to combat premature convergence are: population size, type and probability of crossover, type and probability of mutation, and selection operators (7:106-124).

Population Size. A genetic algorithm population contains a fixed number of strings. This number is called the *population size*. The population size affects the convergence rate of a genetic algorithm by controlling the variety of genes in the population. A smaller population may converge quickly, but usually to a sub-optimum. A larger population converges more slowly, and usually, but not always, finds a better final answer.

Although the basic procedure is simple, the framework of a simple genetic algorithm has been successful in solving a wide range of problems in function optimization and has been extended to other types of problems such as combinatorial optimization (13:165,193).

Why GAs Work. Although there is no complete formal theory to explain the operation of genetic algorithms, several hypotheses which partially explain their power have been advanced (13:51). The point they make is that in each generation, the genetic algorithm combines good partial solutions in the genes of parent chromosomes to find even better solutions in child chromosomes. The following discussion briefly describes these hypotheses. A more formal treatment of genetic algorithm theory is available in Goldberg's text (7).

Schemata and Schema Theorem. A schema is a pattern of values in a gene with the alphabet 1,0,*, where "*" is the "don't care" symbol matching any position. For a chromosome to include a particular schema it must match the schema values. For example, the chromosome (01101) contains 32 schemata. These include: (0**01), (*1****), (01101), (*****), (0*1*1). The *defining length* of a schema is the distance between the outermost non-* symbols. The *order* of a schema is the number of non-* symbols contained in the schema.

Holland's schema theorem was the first rigorous explanation of how a simple genetic algorithm works. The schema can be thought of as representing the partial solutions in chromosomes and Holland concluded that genetic algorithms manipulate schemata when they execute (13:91). In the simple genetic algorithm, individuals reproduce and increase in proportion to their fitness. The schema theorem asserts that because of this, schema associated with individuals with above average fitness tend to increase exponentially, while those schema associated with below average performance tend to occur less often in succeeding generations (7:32-33).

Building Block Hypothesis. Much of the power of genetic algorithms comes from finding good building blocks (7:41). Building blocks are highly fit, low-order schema of short defining length. Because of their short length, these blocks tend to survive, even under the disruption caused by crossover. "In a way, by working with these particular schemata (the building blocks), we have reduced the complexity of our problem; instead of building high-performance strings by trying every conceivable combination, we construct better and better strings from the best partial solutions of past samplings" (7:41).

Together, the schema theorem and the building block hypothesis help to explain why genetic algorithms work. However, the hypothesis does not give a formula for designing genetic algorithms. Instead, most practical information on genetic algorithm design and performance have come from empirical studies.

Permutation Genetic Algorithms

The simple genetic algorithm with a binary coding is appropriate for many unconstrained optimization problems. However, in many *order-based* problems, the solution may be specified by a specific arrangement of items. Examples include scheduling problems and the Traveling Salesman Problem (TSP). Traveling Salesman Problem. The Traveling Salesman Problem is easily stated: a salesman must visit customers in each of n cities without visiting a city twice (13:165-167). The objective is simply to travel the least total distance and return to the starting city. A sample network of cities is shown in Figure 4. The solution to the TSP can be represented as a list of integer numbers with each integer corresponding to a city, and the cities visited in the order of the list. For example, a starting solution for a TSP with six cities could be represented as: A = 1 2 3 4 5 6. This solution visits each city in ascending order from "1" to "6", returning to "1" to complete the tour. All possible solutions to the TSP can be

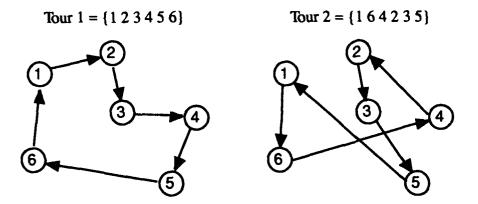


Figure 4. Sample TSP tours

represented as a permutation of the list of integers. A *permutation* is simply an arbitrary reordering of a set of items in a list. One permutation of the 6-city problem is A = 1 6 4 2 3 5. Although the ordering of the cities has changed, the number of cities remains the same and no city values are repeated in the list.

With *n* cities, there are $\frac{(n-1)!}{2}$ possible tours; a TSP with n = 6 has 60 $(\frac{5!}{2})$ possible solutions. For a TSP twice as large (n = 12), the total number of permutations is 239,500,800. The number of possible solutions is obviously incredibly large. Although smaller problems can be solved via deterministic graph search formulations, most large TSP problems are solved (non-optimally) using heuristic techniques which take advantage of distance information between the cities (12).

Blind Traveling Salesman Problem. Note that the TSP being solved by the genetic algorithm is harder to solve than the typical TSP because it does not use any distance

information to solve the problem. "In the blind traveling salesman problem, the salesman has the same objective with the added restriction that he is unaware of the distance he travels until he actually traverses a complete tour" (7:170). However, other ordering problems may not have "distance" information that can be exploited by a heuristic. As pointed out by Whitley, "This is important, because it means the method may be used on sequencing problems where there are no actual distances to measure, but rather only some overall evaluation of the total sequence" (19:137).

Since genetic algorithms had been applied successfully to other optimization problems, it seemed natural to attempt a genetic algorithm solution to the TSP. However, when standard genetic algorithms are applied to the TSP, they have difficulty in finding good feasible solutions (9). The major problem with a standard crossover operator is that most solutions are infeasible (i.e., $1\ 2\ 2\ 4\ 5\ 6$) and the genetic algorithm probably converges on sub-optimal solutions (13:167). Some early attempts used a hybrid technique to combine genetic algorithm crossover with a repair mechanism to make valid tours (9). This approach was somewhat successful, but other approaches have been attempted using order-based crossover operators (13:168-191).

Order-based Crossover Operators. Another approach to solving ordering problems uses a different representation than the standard genetic algorithm. Instead of binary digits, an order-based crossover operator uses a chromosome that directly represents a solution. This chromosome is simply a list of integers; an example is a solution to the six-city TSP problem described previously (164235).

To exploit the information in the chromosomes, an order-based crossover operator, like simple crossover in the standard genetic algorithm, preserves part of the first parent while incorporating information from the second parent. The position of the genes is important. This contrasts with a standard genetic algorithm where their *value* is important. For the benefits of genetic algorithms to be realized for ordering problems, crossover operators must find other ways of combining information from two parents to build offspring with better fitness. Although not strictly following the standard schema theorem for binary-based genetic algorithms, researchers have developed the concept of *ordering schemata* for order-based genetic algorithms (7:175–179). In an ordering problem, the absolute or relative *positions* of the items are important.

Ordering Schema. Ordering schema use a "don't care" symbol "!" to represent unfixed positions in a string. This is different from the "*" symbol presented earlier which represented unfixed values in a string. For example, in a position-based schema, the string (! ! 6 4 ! !) represents cities 6 and 4 in the third and fourth positions, with the "!" position filled arbitrarily from the remaining cities. Possible strings with this position schemata include: (5 1 6 4 2 3), (2 5 6 4 1 3), (3 2 6 4 5 1).

Another ordering schema uses relative ordering. As described by Davis(5:79):

It is important to understand that what is being passed back and forth here is not information of the form "node 3 is in the fifth position." Instead this operator combines information of the form "node 3 comes before node 2 and after node 7." The schemata in a *n* order-based representation can be written in just this way. Let us denote the nodes in a permutation by their indices. The chromosome (516423) contains a number of schemata, including (54), (1423), (563), and (516423).

New crossover operators were developed that could work with ordering problems. Examples include Partially Matched Crossover (PMX) (7:154), position-based crossover (16:343), and edge-recombination (19). These methods attempt to retain the benefits of crossover while maintaining feasible solutions. Two of these operators are described below.

PMX. PMX was developed by Goldberg for use in solving TSPs. He describes the operator using a ten-city problem as an example (7:171). Each city is visited in ascending order; a sample permutation is: (1 2 3 4 5 6 7 8 9 10).

Under PMX, two strings (permutations and their associated alleles) are aligned, and two crossing sites are picked uniformly at random along the strings. These two points define a *matching section* that is used to effect a cross through position-by-position exchange operations.

To see this, consider two strings:

$$A = 9 8 4 5 6 7 1 3 2 10$$

B = 8 7 1 2 3 10 9 5 4 6

PMX proceeds by positionwise exchanges. First, mapping string B to string A, the 5 and the 2, the 3 and the 6, and the 10 and the 7 exchange places. Similarly mapping string A to string B, the 5 and the 2, the 6 and the 3, and the 7 and the 10 exchange places. Following PMX we are left with two offspring, A' and B':

$$A' = 9 \ 8 \ 4 \ 2 \ 3 \ 10 \ 1 \ 6 \ 5 \ 7 \ B' = 8 \ 10 \ 1 \ 5 \ 6 \ 7 \ 9 \ 2 \ 4 \ 3$$

where each string contains ordering information partially determined by each of its parents (7:171).

Position-based Crossover. Another crossover operator, position-based crossover, was developed by Syswerda (16) and is called uniform order-based crossover by Davis (5). This operator attempts to preserve information about the relative ordering of elements in each of the parents. Several random positions are selected from each parent. These positions are inherited by one child. The other positions in this child are inherited in the order they appear in the other parent, skipping over all those included by the first parent. For example, using the strings used to illustrate PMX, first generate a bit string that is the same length as the parents:

A = 98456713210 011001010 B = 87123109546

Next, fill in some of the positions on Child 1 by copying them from Parent A wherever a "1" appears in the binary template.

A list of the elements in Parent A associated with "0" is made, and permuted to appear in the same order as in Parent 2. These permuted elements fill in the gaps in Child 1 to complete crossover. List of elements associated with "0": (9,5,6,1,3,10)

Permuted in order of Parent 2: (1,3,10,9,5,6)

Child
$$1 = 18431079526$$

A similar process is used to create a second child, but with the roles of Parent A and B reversed.

Although PMX and position-based crossover are effective for ordering problems, they process different kinds of ordering schema. PMX tends to respect absolute city position, whereas position-based crossover tends to respect relative city position. Davis explains:

The information that is encoded here, however, is not a fixed value associated with a position on the chromosome. Rather, it is relative orderings of elements on the chromosome. Parent 1 may have a number of elements ordered relatively well. Uniform order-based crossover allows Parent 2 to tell Parent 1 that others of its elements should be ordered differently. The net effect of uniform orderbased crossover is to combine the relative orderings of nodes on the two parent chromosomes in the two children (5:79).

These crossover operators are mentioned because they were used during the implementation of this research. They are effective for problems where ordering is important. PMX is more important where absolute position matters, while position-based crossover is more effective where a relative ordering is more important.

Genetic Algorithms in Scheduling

Many recent papers involving genetic algorithms deal with scheduling problems. Two main approaches to solving scheduling problems with genetic algorithms have been developed: direct chromosome representation and indirect chromosome representation.

Direct Chromosome Representation. Bruns advocates a direct chromosome representation for scheduling problems. In a direct problem representation, the production schedule itself is used as a chromosome. No decoding procedure is therefore necessary. The extended chromosome representation requires the construction of domain-specific recombination operators (4:355) and would therefore be problem-specific. In this case: The sequence of the items within a chromosome is of no importance. To determine the quality of a chromosome any arbitrary evaluation function that has been used in traditional scheduling approaches is applicable without any prior transformation. (4:356)

According to Bruns, this approach should perform better than domain-independent operators, although he admits he cannot rely on any theory (4). It also involves creating custom crossover operators for each problem, which may be difficult. However, this approach has the advantage that the genetic algorithm is allowed to search the entire solution space, not just the ordering of the requests.

Indirect Representation. An indirect representation may also be referred to as a *hybrid* technique (7:202) where part of the problem is solved with the genetic algorithm and the other is solved using a deterministic routine. The first part is a sequencing problem, solved by the genetic algorithm, which orders each job request in a list. The second part is a schedule builder which takes each job, in the order of the list, and attempts to place each job request in a place in the schedule without overlapping another scheduled job. Obviously, those jobs earlier in the list are easier to schedule since more room in the schedule is available.

The schedule builder uses a simple rule to schedule each job; usually, a job is scheduled in the first position which meets the constraints of the problem. Although a schedule builder can use some information to find a good place in the schedule for each job, it is usually best to attempt the genetic algorithm search for good overall schedules by permuting the list of supports (16).

Much research on genetic algorithms in scheduling has been conducted recently in areas such as job shop scheduling and vehicle routing (3:452-459). For application to the satellite range scheduling problem, resource and sequence scheduling are more relevant. Two of these are reviewed by way of example in the next section.

Scheduling Examples

F-14 Test Range Scheduling. To solve a resource scheduling problem involving a test laboratory for F-14 fighter aircraft, Syswerda separated the genetic algorithm from the specific problem (16:332). The list of items to be scheduled is represented as a string of numbers. The genetic algorithm permutes the order of the items in this string to find the best order in which to schedule the items. Then, given the ordered list of items produced by the genetic algorithm, a schedule builder program builds a feasible schedule. The schedule builder is merely a program which attempts to schedule each item in the order presented by the string. The number of items successfully scheduled is returned to the genetic algorithm as a fitness score. This type of approach to scheduling assumes that given a correct ordering of tasks, the schedule builder program can build the best schedule. Syswerda noted:

One thing that is clearly important, especially with regard to the greedy considerations of a single task, is the position of that task in the list. The closer the task is to the front of the list, the greater is its chance that it will be placed into the schedule ... if two tasks both require a scarce resource, the first task in the list may prevent the second from being scheduled, implying that the order of tasks is also important. (16:340)

Syswerda's implementation is interesting because he was able to satisfy many scheduling requirements such as priority of items and user preferences for scheduling days. Such flexibility is important for successful implementation of a scheduling solution.

Coors Scheduling. Much work has been done at Colorado State University on developing genetic algorithms for application to the Traveling Salesman Problem and some scheduling problems. Whitley developed a genetic algorithm called GENITOR (described in Chapter III, and used for this research), and Whitley and Starkweather developed an order-based crossover operator called *genetic edge recombination* for use in solving traveling salesman problems. Along with such theoretical developments, they also applied genetic algorithm solutions to a warehouse/shipping scheduler at Coors (15:74), and a production line scheduler at Hewlett-Packard (18:358-360).

From this research Whitley generalizes the application of genetic algorithms to scheduling problems in general:

... a broad class of scheduling problems can be viewed as sequencing problems. By optimizing the sequence of processes or events that are fed into a simple schedule builder, optimization across the entire problem domain can be achieved. Schedules have not always been viewed this way because there has not existed a general purpose mechanism for optimizing sequences that only requires feedback about the performances of a sample sequence... Thus, a "genetic" approach to scheduling has the potential to produce some very general scheduling techniques, and could be the foundation of a general purpose approach to sequence scheduling. (18:358)

For this research, an approach similar to that used by Syswerda in the F-14 test range scheduling problem seems promising, as it does not require a custom chromosome. Hence, existing genetic algorithm packages, such as GENITOR, may be used for implementation. This approach only requires selecting an ordering crossover operator and constructing a schedule building program to simulate the scheduling operation. As described in the next chapter, such a strategy can be easily formulated.

Other Solution Efforts for Satellite Range Scheduling

Solutions to the satellite range scheduling problem have been studied recently. These include an effort by IBM, and two thesis efforts at the Air Force Institute of Technology (AFIT). These efforts used mixed-integer programming and heuristic approaches to find solutions.

Arbabi's Approach. The first study of automating satellite range scheduling took place during the 1981-84 time period when IBM conducted a study to determine the feasibility of automating satellite range scheduling (1:271-277). Arbabi concluded that a mixed-integer programming approach was not feasible for problems with more than 50 requests. Instead, he developed an approach called Continuous Time Scheduling (CTS). The procedure was not described in detail, as it is apparently proprietary, but it used a heuristic approach. This procedure reportedly scheduled 92% of the requests for one day (1:277).

Gooley's Approach. Gooley used both a mixed-integer programming approach and heuristic scheduling methods (8). The satellite range scheduling problem was successfully formulated as a mixed-integer program (MIP), but the number of integer variables prohibited direct solution. To reduce the number of integer variables, Gooley divided the problem into two mixed-integer programs which scheduled low-altitude satellite supports (the MIP could handle up to 85 low-altitude supports at once). He then used heuristic insertion and interchange techniques to schedule the medium- and high-altitude satellite supports. Using this method, a test set of problems was solved in under 20 minutes with approximately 92% of all requested supports scheduled (8:5-2).

Schalck's Approach. In a follow-on thesis effort, Schalck improved on Gooley's solution by reducing the number of integer variables needed in the MIP formulation. By doing so, one MIP solution scheduled all low-altitude supports for one 24-hour period. Solution times for scheduling the high-altitude satellites in one 24-hour block was too long, and was reduced to about 30 minutes by scheduling high-altitude supports in two 12-hour blocks. The resulting solution scheduled approximately 98% of requested supports (this number is not directly comparable to Gooley's results because of differences in how the support requests were generated).

The mixed-integer programming approach taken by Gooley and Schalck is the best approach for the satellite range scheduling problem when the problem is small enough to be solved by the mixed-integer program, since such an approach should find an optimal solution. However, by decomposing the problem into separate problems, an overall optimal solution may no longer be guaranteed. For example, splitting the requests into two blocks means that support requests near the division point may not be scheduled. A genetic algorithm approach may produce better schedules while meeting the constraints of a short solution time because it would attempt to find a solution to the entire problem at once. Although the genetic algorithm approach is not guaranteed to find an *optimal* solution, it can attempt to find good solutions. Also, a genetic algorithm approach may be more flexible in handling additional constraints for special scheduling requests.

Summary

This chapter summarized the development of genetic algorithms including the standard genetic algorithm and variants. The extension to order-based genetic algorithms allows good solutions to some combinatorial optimization problems, such as the TSP. This in turn allows applications of genetic algorithms to scheduling problems by dividing the algorithm solution between a deterministic schedule builder and a genetic algorithm. Finally, past research efforts in satellite range scheduling were reviewed, with the conclusion that better solutions may be found by scheduling the entire day's schedule in one time block instead of decomposing the problem into smaller time blocks. This may be done with a genetic algorithm based approach.

III. Solution Methodology for Satellite Range Scheduling

This chapter presents a review of the satellite range scheduling (SRS) problem, followed by the development of a genetic algorithm-based scheduling strategy. By formulating the scheduling problem as a sequencing problem rather than as a mathematical program, a straightforward solution by an order-based genetic algorithm is possible. The chapter ends with a review of the implementation of the approach using the genetic algorithm package GENITOR.

Formulation

Definition of the Problem. In the SRS problem, satellite communication supports compete for RTS (remote tracking station) time in a 24 hour schedule. Each support must be scheduled in a restricted time window at certain RTSs due to visibility and scheduling requirements. For the low-altitude satellites, the requested support length fills the entire window. For medium to high altitude satellites, the time window is more flexible, as a tolerance for the beginning of the support is allowed for scheduling. Although each lowaltitude support requires a fixed time window at one RTS and fills a time-window, the scheduling task is eased as most RTSs have two sides (antennas) capable of supporting communications. The RTS sides can support a satellite support simultaneously; if a satellite is visible to an RTS, it can be supported by any one of the available sides at the RTS. In addition, most medium-to-high altitude satellites are visible to more than one RTS. This fact, combined with the more flexible time windows of these satellites, makes scheduling them easier than scheduling low-altitude satellites.

Small Problem. An example set of time windows for a small set of five support requests is shown in Table 1. "Spt" is an arbitrary support number, "Begin" is the starting time for the window, "End" is the ending time for a window, "Length" is the actual service time needed, and "TAT" is the setup time required before a service time can begin. These supports are all serviced by RTS "POGO-A."

This sample time window data can be used to illustrate the satellite range scheduling (SRS) problem. This small problem is simplified since a real day's schedule would include

Sr :	Begin	End	Length	TAT
1	1	13	3	1
2	15	22	3	2
3	7	17	3	1
4	1	10	3	1
5	2	8	3	2

Table 1. Small Schedule Time Windows

over 300 supports, nearly all of which would have alternate windows for scheduling. These alternate windows could come from different antennas or sides at the same RTS, or from different RTSs. In addition to the table, this information is shown graphically in Figure 5.

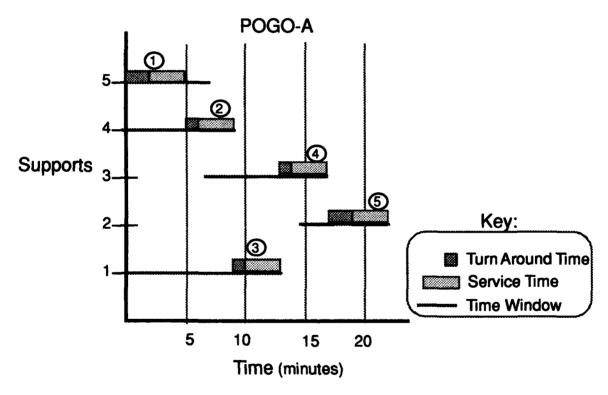


Figure 5. Small Schedule

The horizontal axis shows time in minutes, the vertical axis shows the supports by support number, and the entire chart is for RTS POGO-A. A problem with more than one RTS could be represented by more than one chart. In Figure 5, the *time window* is represented by a thin line, while the required TAT and service time is shown by the boxes (darker pattern for TAT). Such a representation clearly shows a schedule. The schedule for the example successfully schedules all supports with no overlap of the support times and no violation of time windows. The TAT for support 5 is outside the service support window, but this is legal as no communications take place at that time, only set up for a support.

Mized-Integer Programming Approach. Previously, the SRS problem has been formulated as a mixed-integer programming (MIP) problem (Gooley and Schalck). For a problem with a small number of variables, this is the best approach because an optimal solution can be found. But, as the number of supports to be scheduled increases, so does the number of integer variables in the MIP formulation. Because of this increase, an optimal solution may no longer be available in a timely manner. Gooley and Schalck address this problem by decomposing the problem into parts. Gooley combined the use of a MIP formulation with heuristics. Schalck accomplished variable reduction and scheduled the satellites using an MIP for various combinations of satellites and blocks of time. These methods produce feasible schedules which are not necessarily optimal.

A MIP solution includes the starting time for each support and the RTS used for the support. The schedule can then be generated by adding the service times to the starting times. With this type of formulation, the mathematical program must find the actual starting times without violating constraints. A MIP solution of the small problem described in the previous section might begin with support 1 scheduled to start at time 10, support 2 scheduled to start at time 19, and so on, as shown in Figure 5.

Scheduling as a Sequencing Problem. Instead of formulating the satellite range scheduling problem as a mixed-integer program, this research approaches the problem as a sequencing problem where a sequence is defined as an ordered list of items. A solution to a sequencing problem in scheduling determines the "best" order in which to schedule items using simple rules for placing each item in the schedule. For example, in the satellite range scheduling problem, the solution is represented as a sequence of supports. In the small problem, there are five supports. An example ordering is: (5, 4, 1, 3, 2). To build a schedule from this representation, each support is scheduled according to the order it appears in the list. In the example schedule, the first support, 5, would be placed in the first available time segment. The entire schedule is open and support 5 is scheduled at the beginning of its time window from time 2 to time 5. Note that the turn-around time of two minutes is scheduled from time 0 to time 2 which is outside the time window. This is valid since no communications take place during the turn-around time. The schedule is then updated to reflect that time 0 to 5 is no longer available. The second support in the list is 4. Its time window begins at 1, but it cannot be placed there because support 5 has already used minutes 0 to 5. Support 4 is then scheduled from time 6 to time 9 (1 minute TAT; 3 minutes for service). This support remains within its time window which ends at time 9. This procedure repeats for supports (1, 3, 2), and completes a schedule as shown in Table 2.

Spt	Begin	End	Length	TAT
5	2	5	3	2
4	6	9	3	1
1	10	13	3	1
3	14	17	3	1
2	19	22	3	2

Table 2. Small Problem Schedule

The actual start and end times for each support were determined by examining the completed schedule. Thus, all the information needed to define a solution is contained in the sequence of supports, and can be "decoded" by following simple deterministic rules. Obviously, such rules are somewhat arbitrary. The "first available" rule, as shown here, is simple. A rule which attempts to find the "best" place for each particular support would also work, but at the expense of simplicity and possibly time.

An arbitrary ordering of supports is not likely to produce a perfect, or even good, solution which schedules all support requests. However, if a solution exists, an ordering can be found which results in the optimal solution. In many cases, more than one ordering results in the same schedule depending on the interdependence of each support request. To find good schedules, an efficient way of generating better alternative orderings is needed. In this research, a genetic algorithm with order-based crossover operators is used to spawn good orderings of the support requests. Genetic algorithms can quickly search the solution space, defined as the permutations of the list of items. These permutations can then be used to build a schedule as described above, and the number of supports scheduled successfully can be used as a fitness measure for the genetic algorithm search.

Hybrid Approach. As noted in the previous section, a schedule can be built from a sequence of supports using a schedule builder program. This schedule builder is required for building *feasible schedules* given a sequence of support by the genetic algorithm. The overall process is shown in Figure 6. Although some decision information can be incorporated into the schedule builder to improve local search, the genetic algorithm should do most of the exploration of the search space. This approach is based on the assumption that if the supports are entered in the schedule builder in a certain order, the greatest number of supports can be scheduled. It is reasonable to expect that many orderings of the supports may exist which produce the same schedule.

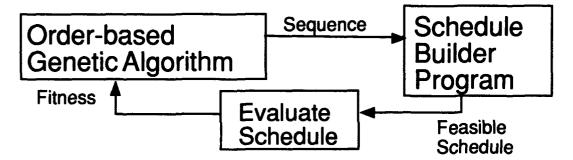


Figure 6. Schedule Builder Flowchart

The nature of the genetic algorithm and schedule builder approach allows other constraints to be imposed on the supports; for example, a request that a support be handled by a certain set of RTSs on a particular day, or that supports must be scheduled exactly one hour apart.

Algorithm Design and Implementation

Overview. Implementation of a hybrid approach requires data preprocessing of satellite requests to generate support request time windows, followed by application of the

GENITOR genetic algorithm. The schedule-building evaluation function is used within GENITOR to produce a given schedule, given a sequence of supports from a GA string.

Assumptions. The research assumes the constraints of the satellite range scheduling problem are those used in the previous efforts by Gooley and Schalck (14:1-7):

- 1. Requested support times are known in advance.
- 2. Time windows for satellite visibility are not flexible. The goal of the SRS problem is to schedule as many supports without having to change the time windows. After the initial schedule is formed, these constraints may be relaxed to schedule unscheduled supports. However, it is desired to minimize the number of such changes.
- 3. Downtimes for RTS maintenance are not included. If known, flexible downtimes could be added as additional supports. Fixed downtimes could be be scheduled before regular supports are scheduled.

Use of these assumptions allows comparison of the genetic algorithm results to previous efforts.

Data Processing. Time windows for each support are processed from the raw satellite visibility and request data in the ASTROS database for a day. The support requests and time windows are represented as lines in a database with the following information:

Support Number: arbitrary support number

- **RTS:** Remote tracking station and antenna side (example: POGO-A for A-side, POGO-B for B-side)
- Beginning of Time Window: in minutes (example: 0200 is 120 minutes; 2400 is 1440 minutes)

End of Time Window: in minutes

Support Length (minutes): actual service time needed

Turnaround Time (TAT): set up time at RTS needed before service time (20 minutes for low-altitude satellites; 15 minutes for medium-high altitude satellites) Satellite Identification: IRON (first four digits; identifies satellite) and revolution num-

ber (last three digits; identifies orbital pass)

An example set of time windows is shown in Table 3. This data is the same as for Table 1, but with fields included for the RTS name and and satellite identification number.

Spt	RTS	Begin	End	Length	TAT	Ident
1	POGO-A	1	13	3	1	2532097
2	POGO-A	15	22	3	2	4774042
3	POGO-A	7	17	3	1	9845009
4	POGO-A	1	10	3	1	3187074
5	POGO-A	2	8	3	2	9757024

Table 3. Small Schedule Time Windows

This time window data is read into data structures at the beginning of a run.

Data Structures. Time window data is read into a structure of arrays. The information can then be used each time a schedule is built from a sequence of supports. These array structures keep track of time window alternatives for each support, and allows the schedule to be filled in as supports are scheduled.

Schedule array. An array named filled (RTS, time) shows the status of each minute (time) for every RTS and antenna combination. A '0' represents an empty minute; a '1' indicates that minute is filled. This array is updated whenever a support is scheduled by setting the minutes used by the support to '1.' Any set up time is also blocked out of the schedule.

Support Information. Three arrays specify the requirements of each support: NumWin(support) is the number of time windows where a support can be potentially supported. Then the support time length requirements and turn around time are given by Length and TAT, respectively. This information is used when running the schedule builder program. For each alternative time window for a support, three arrays give specific information. BVIS and EVIS give the beginning time and ending time, respectively, of each time window. RTS specifies the particular RTS used for the support.

Schedule Builder Program. The schedule builder is implemented as a program in the C language since many genetic algorithms are written in C. It is used as an evaluation function for the genetic algorithm by building a schedule from a sequence of supports generated by the GA, as shown in Figure 6. To construct the schedule builder, an insertion program was written, modeled loosely on those used by Gooley (8), which attempts to schedule each support into an available time window. Many alternative schedules are possible, as the insertion rules for the schedule builder are arbitrary. For example, the simplest implementation attempts to schedule supports to the first available numbered RTS. A schedule could also be built by assigning a support to an available RTS with the most open space left in its schedule. The "first available" approach was chosen for simplicity, and to schedule supports next to each other to utilize the available time at the RTSs. This approach simply schedules satellite communication supports to the first available position in time, and across each RTS, starting with the first RTS.

Each support is scheduled in the order it appears in the ordered list of support requests. The schedule builder tries the first RTS where there is a window for the support. If this fails, it tries the next until it runs out of windows. If the support *is* scheduled, the schedule score is incremented and the space used is blocked out of the schedule. Pseudo-code for the program is as follows:

```
Empty schedule
Attempt to schedule each support
Try each window until succeed or exhaust windows
Try each set of empty spaces until succeed or reach end
Update schedule if support scheduled
Update score if support scheduled
Output final score
```

This can also be represented by a flow diagram, as in Figure 7.

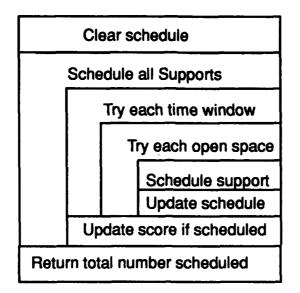


Figure 7. Schedule Builder program flow

This procedure finds the schedule corresponding to a sequence of supports. It is also used to build a schedule given a sequence of supports or to evaluate a given sequence as part of a genetic algorithm implementation.

Genetic Algorithm Implementation. To find "good" support sequences, the schedule builder is integrated into a genetic algorithm as an evaluation function. The genetic algorithm, GENITOR, and associated parameters are described,

GENITOR. Whitley's GENITOR (GENetic ImplemenTOR) code was chosen for use in this research because it has been used with success for other scheduling problems (17). GENITOR includes crossover operators for order-based genetic algorithms (edge recombination, order, and PMX) which have been shown to be useful in solving many scheduling problems (15). GENITOR differs from the standard genetic algorithm as described in Appendix A.

The evaluation function assigns a raw fitness measure to each population member. In GENITOR, the fitness is assigned by an evaluation function written in the C programming language (17). The function input parameter is the chromosome of a population member which is a list of supports. After attempting to schedule each support in the schedule

builder, the function returns a fitness value. In the satellite range scheduling problem, the chosen fitness measure is the number of supports successfully scheduled.

GENITOR Parameter Settings. In GENITOR, the major settings are the crossover operator, population size, selective pressure, and mutation rate (for binary-encoded chromosomes):

Crossover. The crossover operator is important in the quality of the answer. GENITOR includes crossover operators for order-based genetic algorithms as discussed in Chapter II. These include edge recombination, order, and PMX.

Population Size. As explained in Chapter II, the population size affects the convergence rate of a genetic algorithm by controlling the variety of genes in the population. A smaller population may converge quickly, but usually to a sub-optimum. A larger population converges more slowly, and usually, but not always, finds a better final answer. For example, a population size of 30 finds a good answer within 1000 reproductions, but converges to a worse solution.

Selective Pressure. Selective pressure is a selection parameter specific to GENITOR. This parameter controls the rate at which a population converges by giving more reproductive opportunities to higher ranking individuals in the population. In GENITOR, selective pressure is usually set between 1.0 and 2.0. For example, a selective pressure of 1.5 gives the top-ranked individual 1.5 times the chance to reproduce than the median-ranked individual. A lower setting slows down convergence, hopefully allowing more time for the best solution to emerge. A higher setting drives the GA towards the final answer more quickly, but often at the expense of the best solution.

Mutation. It is important to keep a diverse genetic pool as the population converges. Mutation randomly changes part of a chromosome and is usually of less importance than crossover. Mutation was not used in this research because an order-based mutation operator was not present in GENITOR. The addition of such an operator may improve results and is discussed in Chapter V.

Day One Results. Various population sizes and selective pressure settings were tested in GENITOR using the first day of data to determine good parameter settings. In preliminary tests, all order-based crossover operators in GENITOR were tested. Positionbased crossover found the best answer more quickly than the others. Position-based crossover was expected to produce good schedules since it explores the relative position of supports. Since supports are competing for resources, it makes sense that the relative order of supports in a scheduling list is more important than an absolute position in the list. Based on these results, the position operator was chosen as the operator for the test set, although the other order operators would be expected to perform nearly as well. The edge-recombination operator is tailored more towards the TSP and did not do as well as the other operators in this scheduling problem because it stresses adjacency information instead of relative ordering. This agreed with past results using edge-recombination for scheduling problems (15:74).

For Day One data, a graph of performance is shown in Figure 8. As GENITOR executes, at a given interval of reproductions it displays the best individual, the worst individual and the population average. The population has converged when these three numbers are equal, and little improvement can be expected. Convergence indicates that the individuals all have the same fitness score; they are either identical or similar enough to evaluate to the same fitness.

At any time during the run, the GA can be stopped and the best individual found thus far can be chosen as the best schedule found. Usually, the best schedule is found before convergence, but there is no way to know this in advance. In the Day One data, the best schedule was usually found at about 4000 reproductions, while the population did not converge until about 6100 reproductions. A genetic algorithm run could be stopped prior to convergence if time is critical.

Vary Parameters. The runs shown in Figures 9 and 10 display the interplay of exploitation versus exploration in a genetic algorithm solution (7:37). A high selective pressure or a small population size leads to quick convergence of the population. This exploitation of the best individuals found so far often leads to a quick, but less satisfactory

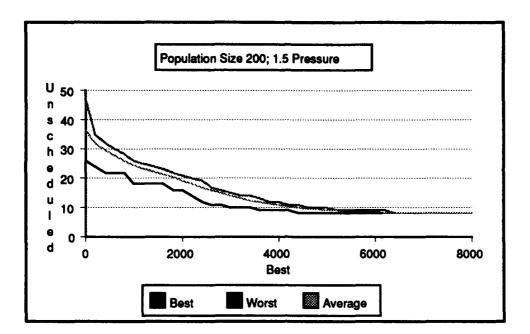


Figure 8. Day 1 Data Convergence

solution. On the other hand, a larger population with lower selective pressure encourages population diversity longer, but at the cost of slower convergence.

Repeated Runs. Since a GA has random components, a set of 10 runs was used for the Day One data in order to initially evaluate the variance of the results. Of 322 supports requested, the number of unscheduled supports was 8 in five of the runs and 9 in the other five runs; the average is 8.5, with a variance of 0.28. For this problem, the difference between runs is minor, and in practice one run should be sufficient for scheduling.

Random Schedule Generation. To give a baseline of the performance of this work, one might ask how well a random permutation of supports would perform. To answer this question, an additional run which generated random permutations of the support list and then used the schedule builder program to create schedules was performed. It found reasonable results within 8,000 reproductions, with a best solution of 23 supports unscheduled out of 322. This is not as effective as the best genetic algorithm schedule of only 8 unscheduled. These results are shown in Figure 11.

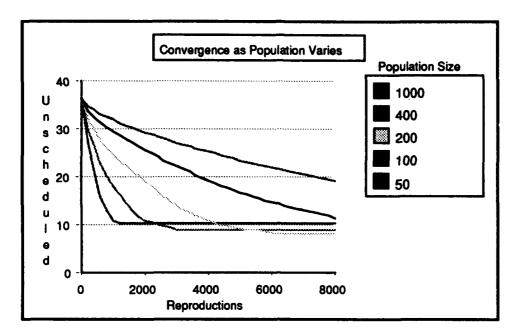


Figure 9. Vary Population Size for Day 1 Data

Overall Implementation. The overall implementation of the genetic algorithmbased approach begins with the processing of raw satellite request and visibility data into a format readable by the GENITOR main program. This processing is handled by Pascal programs written by Gooley and Schalck. After processing, the time window and request data for each support is stored in an array structure which the schedule builder uses to build feasible schedules.

Summary

This chapter describes how the SRS problem may be solved by scheduling items in an order determined by a genetic algorithm. By using this evaluation of each schedule to determine fitness in an order-based genetic algorithm, better schedules can be developed. The genetic algorithm package chosen for this research is GENITOR, because of its success in similar problems. Experiments using the first day of data provide good parameter settings for the GENITOR genetic algorithm in producing 24-hour schedules. The final settings are: position-based crossover, population size of 200, and selective pressure of 1.5. These settings gave a good compromise between the greatest number of supports scheduled and execution time. These settings are used to produce the schedules in Chapter IV.

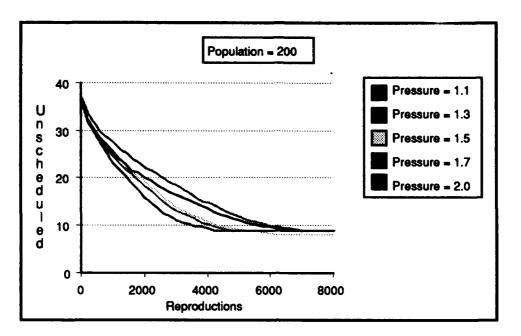


Figure 10. Vary Selective Pressure for Day 1 Data

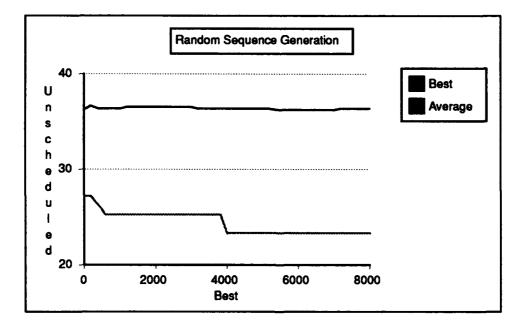


Figure 11. Random Sequence Generation for Day 1 Data

IV. Results

In order to evaluate the performance of a genetic algorithm approach to satellite range scheduling, the schedule builder model described in Chapter III was integrated into the GENITOR (17) genetic algorithm. This genetic algorithm was then tested on real satellite range data. In nearly all cases, the GA-based solution was able to match or exceed previous results that used a mixed integer programming (MIP) approach.

The primary purpose of this research was to evaluate the performance of the GAbased solution in scheduling satellite supports. A secondary purpose was to explore the flexibility of the GA-based approach in handling additional constraints on the scheduling of supports. Such flexibility would be important in implementing a genetic algorithm-based approach in the real world. The additional constraint tested here is the priority of different supports.

The criteria used to judge scheduling results is the number of supports scheduled and the time required for the solution. Supports not scheduled in the initial 24-hour schedule by the range schedulers must be deconflicted through coordination with the MCCs. In the worst case, those supports that cannot be rescheduled by relaxing constraints may have to be canceled. Short solution times are important if different parameters or priorities are to be used by the schedulers. Short solution times are also important for rescheduling on short notice.

Test Set

Scheduling data used in previous research efforts was available and used to test the hybrid GA approach, and to allow a comparison between these results and results of previous efforts. This data is taken from the ASTRO (Automated Scheduling Tools for Range Operations) (1:271) database which is currently used to assist manual scheduling. The primary data is from the seven day period from 12 Jul 92 to 18 Jul 92. This is the same data used by both Gooley and Schalck. Each day includes approximately 300 support requests and their associated request windows. Of these, approximately half are low-altitude satellite support requests. The database contains satellite support requests and the visibility windows of the requests to remote tracking stations for each day. Pascal programs developed by Gooley and modified by Schalck were used to to put the requests in a format suitable for further work. Schalck's processing of the data was used here to allow comparisons to his results. After processing, the final data tables list the visibilities and requests for each day as shown in Table 3 of Chapter III. Processing details are available in Appendix D.

Although each day is scheduled separately, supports are allowed to overlap into the next day if their time window extends into the next day. This is included to match the approach of Schalck. The genetic algorithm implementation schedules across days by keeping track of the overlap into the next day and scheduling these overlaps in the next day's schedule before scheduling regular supports.

The schedule builder code developed in Chapter III serves as the evaluation function in GENITOR. The schedule builder produces a valid schedule from an ordered list of supports for each string in the population. The number of supports successfully scheduled is the fitness of each string. This fitness value is then returned to the genetic algorithm. The fitness is used by GENITOR to rank a particular schedule in relation to others in the population. The more fit strings have a higher probability of being selected to reproduce and exchange their genetic information with other strings through crossover.

Experimental Procedure

Tests were conducted in two major stages. The first day's data was tested to find good parameters for the GA (as shown in Chapter III). These parameters included population size and selective pressure. The primary goal of these parameters is to facilitate achievement of the best solution within the shortest time, where the solutions are scored by the number of supports scheduled. Good general settings were sought which would work well for the data sets. In practice, the general settings would not need to be changed each time a new set of data is introduced. The final settings included position-based crossover, a population size of 200, and a selective pressure of 1.5. These settings appear to be robust in giving a good compromise between the best results and execution time.

These parameters were used to produce schedules for all seven days. Experimental runs were performed on Sun Sparc-10 and Sparc-2 workstations. Solution times for one day of data averaged 10 minutes on a Sparc-10 workstation and approximately 24 minutes on a Sparc-2 workstation. On the Sparc-10, producing 1,000 schedules took approximately one minute.

Week One Results. Once a good set of parameters was found, schedules for all seven days of data were produced. The genetic algorithm (GA) results are from GENITOR with a population of 200, selective pressure of 1.5, and up to 8,000 reproductions. Overall results are shown in Figure 12 and Table 4. Supports whose time windows extend into the next day were allowed to be scheduled into the following day. These overlaps were then blocked out of the available RTS time for the next day. This is done by scheduling these overlap requests before the normal support requests in the following day's schedule. The corresponding results from Schalck's mixed-integer programming (MIP) solutions are included for comparison (14:4-2).

The total number of supports requested and scheduled are shown in Table 4. The difference between the MIP and GA-based results are labeled "GA-MIP."

Day	# Requested	GA# Scheduled	MIP# Scheduled	GA - MIP
1	322	314	312	2
2	302	296	296	0
3	311	307	304	3
4	318	315	311	4
5	305	301	299	2
6	299	292	292	0
7	297	291	291	0

Table 4. Results for Week One Data

The genetic algorithm results compare favorably to the MIP results of Schalck. In every case, the number of total supports the genetic algorithm schedules is at least as good as the number of supports scheduled by the MIP.

Breakdown by Support Type. Results are also broken down by low-altitude and high-altitude supports. As mentioned in Chapter I, low-altitude satellite supports are usually given higher priority in scheduling since they are more difficult to reschedule. Here the GA scheduled more high-altitude satellites than in the MIP solution and scheduled fewer low-altitude satellite supports than the MIP approach. This is due to the fitness of a schedule being calculated as the total number of supports scheduled without reference to the support type. Schalck's solution guaranteed scheduling of the greatest number of lowaltitude supports possible because they were all scheduled before attempting to schedule any high-altitude satellites. This guarantees the scheduling of the greatest number of lowaltitude satellites, but this may not always be necessary. Since set up times for low-altitude supports may be flexible, adjustment of the low-altitude satellite setup times may allow scheduling more supports overall.

Day	# Requested	GA# Scheduled	MIP# Scheduled	GA - MIP
1	153	147	149	-2
2	137	132	134	-2
3	146	143	143	0
4	142	139	140	-1
5	142	139	139	0
6	144	138	138	0
7	142	138	138	0

Table 5. Low-altitude Results for Week One Data

Additional Runs

Since one of the advantages of a hybrid GA approach is flexibility, other constraints on the scheduling of supports should be allowed. The first obvious constraint for inclusion is the scheduling of priorities for different support types.

Day	# Requested	GA# Scheduled	MIP# Scheduled	GA - MIP
1	169	167	163	4
2	165	164	162	2
3	165	164	161	3
4	176	176	171	5
5	163	162	160	2
6	155	154	154	0
7	155	154	153	1

Table 6. High-altitude Results for Week One Data

Support Priority. Low-altitude supports usually take precedence over high-altitude supports since they are more difficult to schedule. To introduce this preference in scheduling, different fitness scores can be assigned to the different types of supports. As an example, low-altitude supports were given a score of two for each one scheduled; the value of high-altitude supports was set to only one. Where a low-altitude support and a highaltitude support compete for the same position in the schedule, the low-altitude support should be scheduled. The overall results are shown in Table 7 and in Figure 13. A breakdown by low-altitude and high-altitude satellites is shown in Table 8 and Table 9.

Day	# Requested	GA# Scheduled	MIP# Scheduled	GA - MIP
1	322	313	312	1
2	302	296	296	0
3	311	307	304	3
4	318	314	311	3
5	305	300	299	1
6	299	293	292	1
7	297	291	291	0

Table 7. Week One Data with Low-Altitude Supports given Priority

Note that more low-altitude supports were accomplished, although at the cost of a tradeoff of some high-altitude supports. Note, however, for Day Two, the GA scheduled two fewer low-altitude supports than the MIP. Upon examination of the schedule, this is due to some high-altitude satellites from Day One scheduled into Day Two.

This type of priority assignment could also be applied to other situations; even so far as to give each satellite a different priority. An array of priority values could be used

Day	# Requested	GA# Scheduled	MIP# Scheduled	GA - MIP
1	153	149	149	0
2	137	132	134	-2
3	146	143	143	0
4	142	139	140	-1
5	142	139	139	0
6	144	138	138	0
7	142	138	138	0

Table 8. Low-altitude Results with Low-Altitude Supports given Priority

Table 9. High-Altitude Results with Low-Altitude Supports given Priority

Day	# Requested	GA# Scheduled	MIP# Scheduled	GA - MIP
1	169	164	163	1
2	165	164	162	2
3	165	164	161	3
4	176	175	171	4
5	163	161	160	1
6	155	155	154	1
7	155	153	153	0

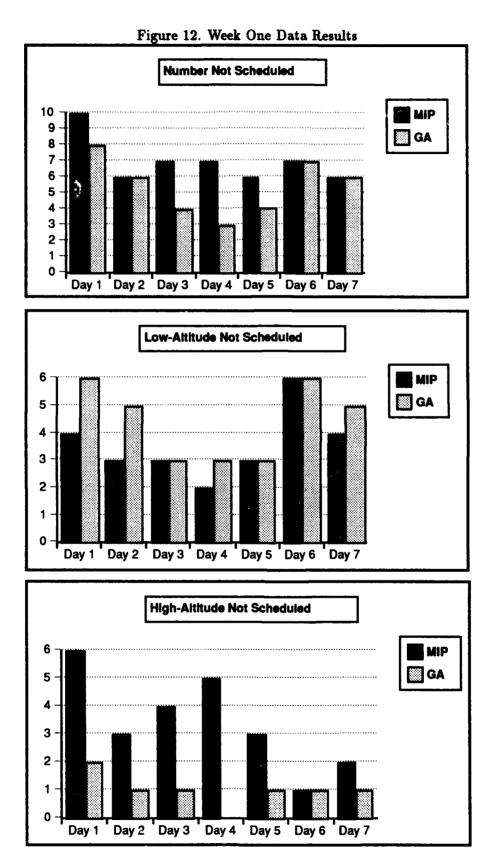
to provide a different score for scheduling each support. A higher score would indicate greater scheduling priority for that support.

Two-Day Scheduling. Even with a priority score given to low-altitude satellite supports, the GA does worse in scheduling low-altitude supports for Days Two and Four. For Day 2, this is caused by the overlap from Day 1. Since each GA only schedules for one day, without regard to the following day, overlaps may reduce the overall number of supports scheduled. The simplest way to overcome this is to schedule two days at once in one block, with priority given to low-altitude supports.

This was attempted with data from Day One and Day Two. Results were encouraging, as the overall number of supports scheduled remains the same, but the number of low-altitude satellite supports scheduled increased to the same level as the MIP solution. In return, the number of high-altitude supports scheduled decreased. Thus, the priority scheme, in conjunction with scheduling for two days, succeeded in matching the low-altitude results of the MIP. Note that the genetic algorithm in this case scheduled two days at once. The solution time was approximately double that for scheduling one day. This indicates that an increase in the number of supports for a day may be handled without much degradation in solution time or quality.

Summary

Results using one week of satellite support data indicates the GA can be successful in scheduling satellite supports. Results match or exceed those returned by the Schalck MIP approach. Solution times are short, and only one computer run is required.



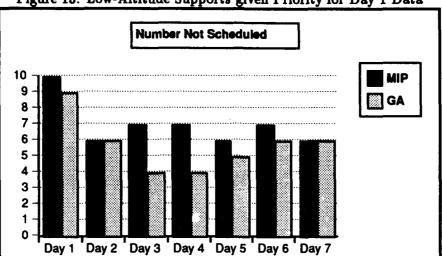
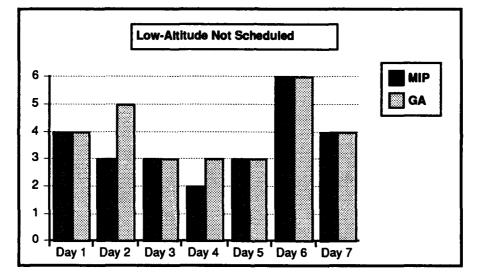
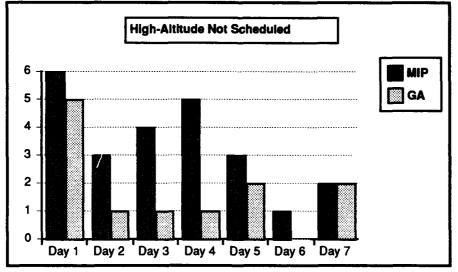


Figure 13. Low-Altitude Supports given Priority for Day 1 Data





V. Conclusions and Recommendations

Conclusions

The satellite range scheduling (SRS) problem involves scheduling over 300 satellite communication supports between Air Force Satellite Control Network (AFSCN) satellites and remote tracking stations (RTS) for a 24-hour period. Scheduling the greatest number of requests minimizes the time needed by schedulers to resolve conflicts in the schedule. The scheduling is currently done manually, with computer aid. Although the current system schedules approximately 95-98% of the requested supports, automating the scheduling process may produce better schedules in less time and use less human resources. The objective of this research was to discover if a genetic algorithm based approach can be effective in automating the scheduling of the requests for a 24-hour period.

The problem can be formulated as a mixed-integer program, but the size of realistic problems precludes producing optimal schedules in a timely manner. By partitioning the original mixed-integer problem into subproblems, Schalck and Gooley found good solutions (14:4-1), but since they do not solve the entire problem at once, there may be room for improvement.

For a test set of seven days of scheduling data, the GA solution matched and often exceeded the results of previous MIP efforts by solving the entire problem at once instead of decomposing the problem into subproblems. The algorithm scheduled over 96% of requested supports. Note, however, these results do not include RTS downtimes. Solution times on Sun Sparc workstations suggest a 24-hour schedule can be produced in under 15 minutes.

These results suggest a genetic algorithm (GA) based method can produce good schedules for the SRS problem. Although the genetic algorithm-based approach described in this research appears to be capable of producing good schedules in a short time, much work remains to be done to implement this automated scheduling approach. Further refinement of the program code may decrease time needed to provide a schedule. More importantly, a fully automated scheduler must also schedule other types of supports which occur in the scheduling process.

Recommendations

Improvements to the algorithm involve further modification and testing of the genetic algorithm with more complicated support requests. The genetic algorithm-based method is easy to implement and modify since all that is required is a schedule builder program to build a schedule one support at a time and to evaluate its fitness score. Since this schedule builder program is easy to test, modifications which allow new constraints can be tested quickly.

Minor Improvements.

Schedule Builder. If the GA-based algorithm is to be implemented, the program code for the schedule builder must be improved. Both readability and code efficiency would lead to better performance and easier modification by users. The program could also be integrated with the processing of satellite request data to make scheduling a one step process for the schedulers, from data processing to final schedule output.

Genetic Algorithm. The genetic algorithm implementation in GENITOR may be improved by adding a mutation operator. Although it does not play as large a role in the genetic search as crossover, mutation acts to keep variety in the population by randomly changing the genes of the population. An order-based mutation operator, such as that advocated by Davis (5:81), would be appropriate for use in this scheduling algorithm. Also, other genetic algorithm packages could be used besides GENITOR. The schedule builder program would simply be used as an evaluation function in the genetic algorithms.

Extensions. To make the produced schedules more realistic, other types of supports should be scheduled. RTS downtimes and scheduling priorities should be determined and included in the schedule. Also, the schedule builder program could be modified to schedule special requests not currently considered by this program. The most important of these would involve scheduling long requests by dividing them into shorter requests which may be scheduled on different RTSs.

Summary. A genetic algorithm based approach offers a way of quickly (minutes) scheduling 24-hour schedules in the satellite range scheduling process. In addition, the general approach could also be applied to other similar scheduling problems where resources must be scheduled. These include ICBM crew scheduling and test range scheduling.

• .

• .

Appendix A. Modified GENITOR Program Code

This appendix describes the GENITOR genetic algorithm, and lists the files modified to use GENITOR with the satellite range scheduling problem. These files include the main program and include files. Although not modified, the source code for the Position crossover operator is included since it is used for the final results.

Genetic Algorithm: GENITOR. The code from an existing genetic algorithm was used for this research. Whitleys GENITOR (GENetic ImplemenTOR) code was chosen because it has been used for similar problems (17). It includes crossover operators for order-based genetic algorithms (edge recombination, order, and PMX). These operators have been shown to be useful in solving many scheduling problems (15).

GENITOR differs from the simple genetic algorithm in two ways. First, instead of creating distinct generations, it creates only one new offspring at a time. Then, rather than replacing the parents with the offspring, GENITOR replaces the lowest ranking member. This approach has been called *steady-state* reproduction (5:35). It works by first selecting two parents from the population and producing two offspring. One offspring is randomly discarded. The other offspring replaces the lowest ranking member of the population. This new string is then ranked in relation to the fitness of the other members of the population, and inserted into the population. The new member can then compete for reproductive opportunities.

The second difference from the standard genetic algorithm is the way fitness is measured. Instead of relying on raw fitness measures, GENITOR uses relative ranking of population members to determine reproductive opportunities. Ranking prevents scaling problems associated with raw fitness values. Scaling problems occur when one individual is so much better that the others that it dominates the population too soon, causing permature convergence of the algorithm. To determine the reproductive opportunities of population members, GENITOR uses linear bias scaling. This linear bias is usually set between 1.0 and 2.0. For example, a bias of 1.5 means the top-ranked individual has 1.5 times the chance of the median-ranked population member to be selected to reproduce (15). Modified Main program Code. The following listing is the main program code for GENITOR, modified for use with the satellite range scheduling problem. At the beginning of a run, this code reads in satellite time window data and overlap data. During execuation of the program, it calls tsp eval.c to evaluate schedules. Finally, it stores the best sequence in the file tourdata.

This main program is modified for use with the Position crossover operator.

```
/* MAIN Program - 31 Jan 94 */
  - 1 Feb, 31 Jan: read/write to files */
/+
/*
  - 28 Jan: add in pre/post processing for overlap */
/*************
                                                     **/
/*
                                                      */
/*
  Copyright (c) 1990
                                                      */
/* Darrell L. Whitley
                                                      */
/*
                                                      */
  Computer Science Department
/*
                                                      */
   Colorado State University
/*
                                                      */
/* Permission is hereby granted to copy all or any part of
                                                      */
/* this program for free distribution. The author's name
                                                      */
/*
   and this copyright notice must be included in any copy.
                                                      */
/*
                                                      */
main_pos.c
 ******
               *****************
#include <stdio.h>
#include <ctype.h>
#include <string.h>
#include "ga_random.h"
#include "gene.h"
#include "ga_global.h"
#include "ga_params.h"
#include "ga_pool.h"
#include "ga_selection.h"
#include "ga_status.h"
#include "ga_signals.h"
#include "op_position.h"
Declare evaluation function which tells you
 how "good" a particular gene's solution is.
 NOTE: the input parameters for the eval
      function should always be a gene and
  the gene's length
 #include "eval_pos.h"
int
main (argc, argv)
int
    argc;
char *argv[];
£
 int
           i,j;
 GENEPTR
          mom, dad, child;
 FILE
          *fp;
         **coord_array;
 int
 CITY_DATA *city_table;
 /* int
            num_diffs; */
/* Character names for input/output */
```

```
int dayNum;
static char overlap[10] = {"overlap0"};
static char tourdata[10] = {"tourdata0"};
static char datafile[10] = {"datafile0"};
/*************
Setup signal handlers.
setup_signal();
Set the global parameters according to command line arguments.
argc--;
argv++;
parse_command_line (argc, argv);
/**************
Print Parameter Values
*****************
fprintf (stdout, "\n");
print_params(stdout);
fprintf (stdout, "\n");
/********************
Input the day number
and concat to file names
printf("What is the day number (1-7)?");
scanf ("%u", &dayNum);
printf("Day Number is %u\n",dayNum);
overlap[7] += dayMum - 1;
                       /* Adds day number to filenames */
tourdata[8] += dayNum;
datafile[8] += dayNum;
printf("%s\n%s\n%s\n",tourdata,datafile,overlap);
/******************************
Seed the Random Number Generator
srandom(RandomSeed);
Allocate a genetic pool referenced by the global, Pool
if ( !(Pool = get_pool(PoolSize, StringLength)) )
fatal_error(NULL);
/********
                    *******
Read in a description of the points to be toured and
create a representation of the distance between them.
**********
/* make_dist_array (ModeFile, StringLength); */
Assign schedule parameters to the program
 /* #include "data.h" */
Read in data
 read_data (datafile,Pool->string_length); /* set = string length */
 read_overlap_data (overlap);
                                   /* 1/28: read overlap data */
/* for (i=1; i<=18; i++)
    printf( " gRTS[%d]=%d, gBeginVis[%d]=%d, gEndVis[%d]=%d,</pre>
gReqLength[%d]=%d, gTurnAroundTime[%d]=%d \n",i, gRTS[i],i,gBeginVis[i],
 i,gEndVis[i], i, gReqLength[i],i,gTurnAroundTime[i] ); */
/* printf("Initialize genetic pool\n"); */
Initialize the genetic pool with data.
          ************************
init_pool (SeedPool, Pool, 0, Pool->size, tsp_eval);
```

```
/* printf("Sort Initial genetic pool\n"); */
/***********************
Sort the initial genetic pool data.
sort_pool (Pool);
/* printf("Allocate temporary storage for parents of reproduction\n"); */
Allocate temporary storage for parents of reproduction, and for child.
mom = get_gene (Pool->string_length);
dad
    = get_gene (Pool->string_length);
child = get_gene (Pool->string_length);
/* printf("Allocate a table to be used with the Order1 (Davis) Operator\n"); */
Allocate a table to be used with the Order1 (Davis) Operator
city_table = get_city_table (Pool->string_length);
/********
Optimize !
*********/
for (/* CurrentGeneration already set :
   either intialized to 0 in its declaration OR initialized by a restart of a previous experiment */;
   CurrentGeneration < NumberTrials;
   CurrentGeneration++)
Choose two genes for reproduction.
get_parents(mom, dad, Pool, linear, SelectionBias);
  Call Order operator to create a child
**********
   /* printf("Start Order operator\n"); */
if (bitgen() )
/* printf("After bitgen = 0\n"); */
position (mom->string, dad->string, child->string, Pool->string_length, city_table);
else
/* printf("After bitgen = 1\n"); */
position (dad->string, mom->string, child->string, Pool->string_length, city_table);
/* printf("So, kid, how good are you?\n"); */
     *******
So, kid, how good are you?
******
   child->worth = tsp_eval (child->string, Pool->string_length);
/* printf("Insert new gene into population according to its worth\n"); */
   /*********
            *****
                 *******
Insert new gene into population according to its worth
insert_gene (child, Pool);
   If the StatusInterval parameter was set and this is the appropriate
time, print the population best, worst, mean, and average to stdout
      ********
              if (StatusInterval & !(CurrentGeneration % StatusInterval))
   show_progress (stdout, Pool, CurrentGeneration);
   If the DumpInterval parameter was set and this is the appropriate
time, save the population and key parameters to disk for later
reference (or to restart execution later.
```

```
*************************
                                                    ********/
    if (DumpInterval && !(CurrentGeneration % DumpInterval))
       dump_status(Pool, DumpBase);
}
/************
 Summarize Results
  final_pool (FinalPool, Pool, CurrentGeneration);
  fprintf (stdout, "\n");
  printf("Right before print_pool\n");
  print_pool (stdout, Pool, 0, 1);
  printf ("After print_pool, before tourdata\n");
  /*************
 Print out tourdata to file
if (fp = fopen(tourdata, "w"))
Ł
              printf("Inside print tourdata\n");
  for (i=0;i < Pool->string_length; i++)
Ł
    fprintf(fp,"%d ",Pool->data[1].string[i]);
3
printf("Its past!\n");
              fclose(fp);
}
else
printf("Cannot print tourdata file");
 }
}
```

Ga global.h Include files. This file contains the global variables used in GENITOR.

The global variables used to store time window information for each support are included here as global variables and used in the tsp eval function. Another related file called ga global external.h is also listed.

```
/*ga_global.h */
/* mod 9 Feb: add temp LF Num for Priority */
/* modified 8 Feb for downtimes; 28 Jan for overlap */
/*
/* Copyright (c) 1990
                                        */
                                         */
/* Darrell L. Whitley
                                         */
/* Computer Science Department
                                         */
/* Colorado State University
                                         */
/*
                                         */
/* Permission is hereby granted to copy all or any part of */
/* this program for free distribution. The authors name */
/*
  and this copyright notice must be included in any copy.
                                        */
/*
                                         */
* These REQUIRED and OPTIONAL parameters can be chosen by the user. *
/* REQUIRED */
     PoolSize;
int
```

```
int
        StringLongth;
 int
        JumberTrials;
char
        ModeFile[80];
                           /* contains coordinates of tsp nodes */
/* OPTIONAL */
long Random
        RandomSeed;
float
        SelectionBias;
        MutateRate = 0.0;
 float
 int
        StatusInterval = 0; /* dump status every nth generation */
 int
        DumpInterval = 0;
                            /* save state of every nth generation */
char
                            /* file containing init data */
        SeedPool[80];
        FinalPool[80];
char
                            /* file containing final results */
 char
        DumpBase[80];
                           /* basename of file(s) into
                              which to dump population */
        NumPop;
 int
                            /* Number of Subpopulations */
 int
        SwapInterval;
                            /* Trials between Swapping of Subpopulations */
 int
        SwapNumber;
                            /* Humber of Strings Swapped between Subpops */
 int
        Experiments:
                            /* Experiments must be used in main() */
float
                            /* Cutoff value for a given experiment */
        CutOff;
int SequenceFlag = 0;
                           /* if set to 1 in main, insert_unique_gene() */
    /* will use different criteria for sameness */
 int
        CurrentGeneration = 0;
POOLPTR Pool;
/* Schedule Builder Global Variables */
/*
     [spt][win] */
int
       gRTS[700][20];
       gBeginVis[700][20]:
int
int
       gEndVis[700][20];
int
       gReqLength[700];
       gTurnAroundTime [700];
int
int
       gBeginSched[700];
int
       gEndSched[700];
int
       gNumWin[700];
int
       gOverlapBegin[20];
       gOverlapEnd[20];
int
int gDownNum;
int gDownRTS[100];
int gDownBegin[100];
int gDownEnd[100];
int gNumLF;
/* ga_global_extern.h */
/* mod 9 Feb: add temp NumLF for priority */
/* mod 8 Feb for downtime */
/* modified 28 Jan 94 to add overlap variables */
/*
                                                           */
/* Copyright (c) 1990
                                                           */
/* Darrell L. Whitley
                                                           */
/* Computer Science Department
                                                           */
/* Colorado State University
                                                           */
/*
                                                           */
/* Permission is hereby granted to copy all or any part of */
/*
   this program for free distribution. The authors name
                                                          */
/*
   and this copyright notice must be included in any copy.
                                                          */
/*
                                                           */
```

```
52
```

/**************************************	***************************************
extern int	PoolSize;
extern int	StringLength;
extern int	JumberTrials;
extern char	NodeFile[];
extern long	RandomSeed;
extern float	SelectionBias;
extern float	MutateRate;
extern int	StatusInterval;
extern int	DumpInterval;
extern char	DumpBase[];
extern char	SeedPool[];
extern char	FinalPool[];
extern int	CurrentGeneration;
extern int	SumPop;
extern int extern int	SwapInterval;
extern int	SwapIumber;
extern float	Experiments; CutOff:
extern int	SequenceFlag;
-	

/	······································
extern POOLPTR	-
extern POOLPTR	-
extern POOLPTR	Pool; ilder Global Variables */ gRTS[700][20];
extern POOLPTR /* Schedule Bu	Pool; ilder Global Variables */ gRTS[700][20]; gBeginVis[700][20];
extern POOLPTR /* Schedule Bu extern int	Pool; ilder Global Variables */ gRTS[700][20];
extern POOLPTR /* Schedule Bu extern int extern int	Pool; ilder Global Variables */ gRTS[700][20]; gBeginVis[700][20]; gEndVis[700][20];
extern POOLPTR /* Schedule Bu extern int extern int extern int	Pool; ilder Global Variables */ gRTS[700][20]; gBeginVis[700][20]; gEndVis[700][20]; gReqLength[700];
extern POOLPTR /* Schedule Bu extern int extern int extern int extern int extern int	Pool; ilder Global Variables */ gRTS[700][20]; gBeginVis[700][20]; gEndVis[700][20]; gReqLength[700]; gTurnAroundTime[700];
extern POOLPTR /* Schedule Bu extern int extern int extern int extern int extern int extern int	Pool; ilder Global Variables */ gRTS[700][20]; gBeginVis[700][20]; gEndVis[700][20]; gReqLength[700]; gTurnAroundTime[700]; gBeginSched[700];
extern POOLPTR /* Schedule Bu extern int extern int extern int extern int extern int extern int extern int	Pool; ilder Global Variables */ gRTS[700][20]; gBeginVis[700][20]; gEndVis[700][20]; gReqLength[700]; gTurnAroundTime[700]; gBeginSched[700]; gEndSched[700];
extern POOLPTR /* Schedule Bu extern int extern int extern int extern int extern int extern int extern int extern int extern int	Pool; ilder Global Variables */ gRTS[700][20]; gBeginVis[700][20]; gEndVis[700][20]; gReqLength[700]; gTurnAroundTime[700]; gBeginSched[700]; gEndSched[700]; gWumWin[700];
extern POOLPTR /* Schedule Bu extern int extern int extern int extern int extern int extern int extern int	<pre>Pool; ilder Global Variables */ gRTS[700][20]; gBeginVis[700][20]; gEndVis[700][20]; gReqLength[700]; gTurnAroundTime[700]; gBeginSched[700]; gEndSched[700]; gWunWin[700]; gOverlapBegin[20];</pre>
extern POOLPTR /* Schedule Bu extern int extern int	<pre>Pool; ilder Global Variables */ gRTS[700][20]; gBeginVis[700][20]; gEndVis[700][20]; gReqLength[700]; gTurnAroundTime[700]; gBeginSched[700]; gEndSched[700]; gHumWin[700]; gOverlapBegin[20]; gOverlapEnd[20];</pre>
extern POOLPTR /* Schedule Bu extern int extern int	. Pool; ilder Global Variables */ gRTS[700][20]; gBeginVis[700][20]; gEndVis[700][20]; gReqLength[700]; gTurnAroundTime[700]; gBeginSched[700]; gEndSched[700]; gWumWin[700]; gOverlapBegin[20]; gOverlapEnd[20]; mWum;
extern POOLPTR /* Schedule Bu extern int extern int gDom	<pre>Pool; ilder Global Variables */ gRTS[700][20]; gBeginVis[700][20]; gEndVis[700][20]; gReqLength[700]; gTurnAroundTime[700]; gBeginSched[700]; gEndSched[700]; gHumWin[700]; gOverlapBegin[20]; gOverlapEnd[20]; mNum; mRTS[100];</pre>
extern POOLPTR /* Schedule Bu extern int extern int goom extern int goom	<pre>Pool; ilder Global Variables */ gRTS[700][20]; gBeginVis[700][20]; gEndVis[700][20]; gReqLength[700]; gTurnAroundTime[700]; gBeginSched[700]; gEndSched[700]; gEndSched[700]; gBumWin[700]; gOverlapBegin[20]; gOverlapEnd[20]; mHum; mRTS[100]; mBegin[100];</pre>
extern POOLPTR /* Schedule Bu extern int extern int gDom	<pre>Pool; ilder Global Variables */ gRTS[700][20]; gBeginVis[700][20]; gEndVis[700][20]; gReqLength[700]; gTurnAroundTime[700]; gBeginSched[700]; gEndSched[700]; gEndSched[700]; gBumWin[700]; gOverlapBegin[20]; gOverlapEnd[20]; mHum; mRTS[100]; mBegin[100];</pre>

Position Crossover Operator

This program code is unchanged, but included since it was used in GENITOR to test the data.

```
*/
/***
            /****
          *********************
 * This is the position operator developed by Syswerda *
 *
   (The Genetic Algorithms Handbook, L Davis, ed) and *
 ۰
   implemented by Susan McDaniel at Colorado State *
 *
    University.
                       *
 ۰
 * To use this program include the following lines in *
 * the main_tsp.c file: *
 *
 ÷
   This call should only happen once: *
 *
   get_city_table (Pool->string_length); *
 * This code should be in a loop so it is called once *
 * for each recombination: *
* if ( bitgen () ) *
      position (mom->string, dad->string, child->string, *
                 Pool->string_length, city_table); *
 *
        else *
      position (dad->string, mom->string, child->string, *
 *
 *
      Pool->string_length, city_table); *
 ****
 * This operator attempts to preserve position *
 * information during the recombination process. *
* Several random locations in the tour are selected
* along with one of the parents and the cities in *
 * those positions are inherited from that parent.
                                                        *
 * The remaining cities are inherited in the order in *
 * which they appear in the unselected parent skipping *
 * over all cities which have already been included in *
 * the offspring. *
 * Example- Position based crossover: *
 * Parent 1: a b c d e f g h i j *
 * Cross Pts:
                                    (Parent 1 selected) *
 * Parent 2: cfajhdigbe *
 * Offspring: a b c j h f d g i e *
 * The cities b, c , f and i are inherited from *
* Parent 1 in positions 2, 3, 6 and 9 respectively.
 * The remaining cities are inherited from Parent 2 as *
* follows: Off[1] = P2[3] since P2[1] and P2[2] have *
 * already been included in the child tour. Then
 * going through Parent 2 in order, Off[4] = P2[4], *
* Off[5] = P2[5], Off[7] = P2[6], Off[8] = P2[8] and *
 * 0ff[10] = P2[10]. *
 *************
                    #include <stdio.h>
#include <malloc.h>
#include "ga_random.h"
#include "gene.h"
#include "op_position.h"
* FUNCTION: get_city_table
 * DESCRIPTION: allocates space for the city data table
 * the city data table contains the position
 * of each city in each parent and also
 * has a field which is set when a city has
```

```
54
```

```
* been included in the offspring tour.
 * IMPUT PARAMETERS: string length
 * RETURN VALUE: the address of the city table
 CITY_DATA +
get_city_table (length)
int length;
£
   CITY_DATA *city_table;
/* malloc one extra location so that cities 1-N can be accessed directly. location 0 will not be used */
   if (!(city_table = (CITY_DATA *) malloc ((length + 1)*sizeof (CITY_DATA))))
      printf ("get_city_table: Malloc failure. \n");
   return (city_table);
}
/******
        * FUNCTION: position
 * DESCRIPTION: performs position crossover operation
 * IMPUT PARAMETERS: two parent gene strings
       space for child string
       the length of the two strings
       the city table address
 * RETURN VALUE:
 void position (dad, mom, kid, length, city_table)
GENE_DATA dad[], mom[], kid[];
int length;
CITY_DATA *city_table;
£
   int num_positions;
   int i, pos, dad_index, kid_index;
   for (i=1; i<=length; i++) { /* initialize city table */</pre>
      city_table[i].used = 0;
   }
   /* select #positions that will be inherited directly from parent */
   num_positions = randomain (2*length/3, length/3);
   for (i=0; i<num_positions; i++) {</pre>
      pos = randomain (length - 1, 0); /* select position randomly */
      kid[pos] = mom[pos]; /* transfer cities to child */
      city_table[mom[pos]].used = 1; /* mark city used */
   }
   dad_index = 0;
kid_index = 0;
   while (kid_index < length) {</pre>
      if (!city_table[mom[kid_index]].used) { /* next position in kid filled*/
  if (!city_table[dad[dad_index]].used) { /*next city in dad not used*/
    kid[kid_index] = dad[dad_index]; /* inherit from dad */
    dad_index ++; /* increment indexes */
    kid_index ++;
 }
```

```
else { /* next city in dad has been used */
    dad_index ++; /* so increment dad index */
    } /* end else */
    } /* end if */
    else { /* next position in kid is filled */
    kid_index ++; /* so increment kid index */
    } /* end else */
    } /* end while */
}
```

Appendix B. Evaluation Function Code

The tsp eval program is a procedure which builds and evaluates a schedule. It takes as input an ordered list of supports to schedule and then, in the order of the list, attempts to schedule each support into a schedule. The function returns a fitness value to the main genetic algorithm.

The schedule builder code includes three procedures: ScheduleWindow, Schedule-Support, and tsp eval. Tsp eval is the main procedure. It is called from GENITOR, and returns the evaluation of a schedule. Tsp eval first clears a schedule, schedules any overlaps from the previous day, then attempts to schedule each support in the order given in the sequence passed to it from GENITOR. ScheduleSupport attempts to schedule each support in the available time windows. Each time window is attempted by calling the ScheduleWindow routine. If a support is scheduled in a window, the space taken by the support is blocked out of the schedule. Then ScheduleSupport returns a fitness value to tsp eval. This fitness value is normally one, but can differ if a priority scheme is used to weight the scheduling of different types of supports. After all supports have been scheduled (or failed to be scheduled), tsp eval returns a fitness score based on the number of supports scheduled to the main GENITOR function.

```
/* SCHEDULE BUILDER PROGRAM */
/* 31 Jan 94: auto file generation */
/* 28 Jan 94 Eval_tsp.c -- position crossover */
/* 28 Jan: add overlap variables */
/* This modification allows supports to be scheduled past the end of the day */
/* if allowed by the supports windows */
/* do this by increasing filled array from 1470 to 2100, and take off end */
/* conditions
#include <stdio.h>
#include <math.h>
#include "gene.h"
#include "op_position.h"
#include "ga_global_extern.h"
int
ScheduleWindow (spt, win, filled)
int
      spt;
int
      win;
int filled[21][2100];
     /* Declare Variables */
     int
          bVis, /* solution /* solution /* ending visibility */
                           /* gBeginVis - gTurnAroundTime : 0 if latter <0 */</pre>
    eVis,
                        /* RegLength + TurnAroundTime */
          rLength.
         endSpt,
                     /* Flag for Support either scheduled or run out */
```

```
/* Counter for support index being scheduled */
        j,
1,
                          /* Counter for filling in scheduled blocks */
                 /*
                                     Index for support */
                    spt:
                          /* current time interval */
         minute,
                     /* length so far */
        length,
         scheduled,
                   /* binary: 1 or scheduled; 0 for not schedule */
         beginSched,
         endSched,
         start,
                     /* current start of support */
tat,
             /* turn around time */
                     /* RTS number -- passed from global gRTS[support] */
         rts;
                       /* scheduledPtr : indicates if scheduled or not */
     /**** Initialize Variables *****/
     scheduled = 0;
    length = 0;
     endSpt =0; /* initialize flag; end when it = 1*/
    /* Initialize Starting Point, Required Length with TAT */
    rts = gRTS[spt][win];
    tat = gTurnAroundTime[spt];
    bVis = gBeginVis[spt][win] - gTurnAroundTime[spt] + 20 ;
     eVis = gEndVis[spt][win] + 20;
/* add twenty for overlap area */
     if (bVis < 0)
£
           tat = bVis + tat;
  bVis = 0;
                   /**** reduces tat if near boundary at start ****/
}
    rLength = gReqLength[spt] + tat;
    minute= bVis + 1;
                       /* my blocks are filled to the right of position */
     start = minute;
                                 /* Starting position of current attempt */
     while (endSpt != 1)
                              /* schedule until successful or end */
     Ł
          if (filled[rts][minute] == 0)
                                           /* if current space free */
                                         /* increment length if space free */
               length += 1;
          else
               length = 0; /* reset to zero if space is filled */
     tat = gTurnAroundTime[spt];
         rLength = gReqLength[spt] + tat;
               start = minute+1;
          }
          if (length == rLength)
                                  /* then schedule the support here */
               1
                   beginSched = minute - gReqLength[spt] - 20;
   endSched = minute - 20;
                   scheduled = 1;
                   for (l = start;l <= minute;l++)</pre>
                      filled[rts][1] = 1;
  }
          if ( (minute == eVis) || (scheduled == 1) )
             endSpt = 1;
          minute += 1;
     } /* End Support */
      /* fprintf (stdout, "scheduled = %d\n", scheduled); */
      /*if(scheduled == 1)
        printf("SCHEDULE Support %d in window %d at RTS %d, begin: %d
   end: <sup>%</sup>d\n",
                    spt, win, rts, beginSched, endSched );
     else
        printf("
                     Support %d not scheduled in window %d\n", spt, win); */
```

return(scheduled);
} /* end ScheduleWindow */

•

```
int
ScheduleSupport(spt,filled)
int spt;
int filled[21][2100];
int i, win, flag;
            /* flag = 1 if scheduled , 0 ow */
int sched;
sched = 0;
                   /* counter for number of windows tried */
win = 1;
flag = 0;
while (flag != 1)
  Ł
     sched = ScheduleWindow(spt,win,filled);
    /* printf ("sched=%d, gWumWin[%d]=%d\n", sched,spt,gWumWin[spt]); */
if ( (sched == 1) || (win == gWumWin[spt]) )
flag =1;
    win +=1;
/* if (sched==0)
 printf("*** Support %d NOT SCHEDULED!***\n",spt); */
return(sched);
float
tsp_eval (order, num_supports)
GENE_DATA order[];
         num_supports;
int
£
     /* Declare Variables */
     int
        filled[21][2100],
                           /* time units in schedule are filled[rts][time] */
                         /* Counter for support index being scheduled */
         j,
        ž,
                       /* Index for support */
        spt,
        numScheduled,
                         /* number of supports scheduled */
                          /* counter for current time */
         minute,
        rtsNum;
                         /* max number of RTSs */
     /**** Initialize Variables *****/
    rtsNum = 20; /* must be one less than max cause of arrays 0-19 */
     /* printf("rtsNum = %d\n", rtsNum); */
    for (k = 1; k \leq rtsNum; k++)
        for (minute = 0;minute <= 2100;minute++)</pre>
        £
            filled[k][minute] = 0;
           /* printf("filled[%d][%d]= %d\n",k,minute,filled[k][minute]); */
        }
     /* Schedule the overlaps first, before the real supports */
     for (k = 1; k <= rtsHum; k++)
 for (minute = gOverlapBegin[k]; minute <= gOverlapEnd[k]; minute++)
     filled[k][minute] = 1;
     numScheduled = 0;
     /* Schedule each Support in Priority Order */
                                      /* Try to schedule support */
     for (j = 0; j< num_supports; j++)</pre>
     £
          spt = order[j];
```

```
/* printf( "Start SS Order[%] = %d\n",j, order[j]); */
         numScheduled += ScheduleSupport(spt, filled);
           /* end num_supports */
    /* fprintf (stdout, "num_supports = %d\n", num_supports); */
   /* The GA wants to MINimize something... so do Max poss - actual */
      /* printf( "\n Wumber Scheduled = %d\n*********************\n\n\n",
numScheduled); */
    return(num_supports - numScheduled);
}
     /* End Eval */
FUNCTION: read_data
 * DESCRIBE: read in data
 * INPUT PARAMETERS: filename of list of satellite support data
 * RETURN VALUE: none
 * CALLS: get_coord_array
         read_coords
get_2D_dist_array
 *
 *
          calc_distances
 *
 ******
                      void
read_data (coord_file, num_supports)
char
              coord_file[];
int
              num_supports;
Ł
  int i;
  int spt;
  int rts;
int bVis, eVis, rLength, tat;
  FILE *fp;
  printf("coord_file is: %s", coord_file);
  for (i=1; i<=num_supports; i++)</pre>
£
gNumWin[i] = 0;
                     /* set number of windows [spt] = 0 */
/* printf("gNumWin[%d]=%d\n",i,gNumWin[i]); */
  if (fp = fopen(coord_file, "r"))
£
         while (fscanf(fp, "%d %d %d %d %d %d %d", &spt, &rts, &bVis,
 &eVis,&rLength, &tat ) !=EOF)
         Ł
              gNumWin[spt] +=1;
              gRTS[spt][gNumWin[spt]] = rts;
              gBeginVis[spt][gHumWin[spt]] = bVis; /* add for overlap */
              gEndVis[spt][gNumWin[spt]] = eVis; /* add for overlap */
              gReqLength[spt] = rLength;
              gTurnAroundTime[spt] = tat;
              /* printf ("spt = %d, rts= %d, Bvis= %d, Evis = %d, ReqLen = %d,
 TAT = %d\n", spt, gRTS[spt][gHumWin[spt]], gBeginVis[spt][gHumWin[spt]],
gEndVis[spt][gHumWin[spt]], gReqLength[spt], gTurnAroundTime[spt] ); */
      /* printf("-----gHumWin[%d] = %d\n", spt,gHumWin[spt]); */
 fclose (fp);
  /* for (i=1; i<= num_supports;i++)</pre>
printf("gWumWin[%d] = %d\n",i,gWumWin[i]); */
3
 else
```

```
61
```

```
fatal_error ("Cannot read input city file datafile");
}
FUNCTION: read_overlap_data
 DESCRIBE: read in data
* IMPUT PARAMETERS: filename of list of satellite support overlap data
* RETURN VALUE: none
# CALLS: get_coord_array
        read_coords
get_2D_dist_array
*
*
*
        calc_distances
                       *****
void
read_overlap_data (coord_file)
char
             coord_file[];
£
 int i;
int rts;
int bVis, eVis;
 FILE *fp;
 if (fp = fopen(coord_file, "r"))
£
       while (fscanf(fp, "%d %d %d", &rts, &bVis, &eVis ) !=EOF)
       £
           gOverlapBegin[rts] = bVis;
     gOverlapEnd[rts] = eVis;
            /* printf("g0verlapBegin[%d] = %d\t g0verlapEnd[%d] = %d\n", rts,
gOverlapBegin[rts], rts, gOverlapEnd[rts]); */
           /* printf ("spt = %d, rts= %d, Bvis= %d, Evis = %d, ReqLen = %d,
fclose (fp);
 /* for (i=1; i<= num_supports;i++)</pre>
printf("gBumWin[%d] = %d\n",i,gBumWin[i]); */
}
alsa
fatal_error ("Cannot read input city file overlap");
3
```

Appendix C. Schedule Builder

This program code is a modification of the Evaluation Function code in Appendix B. Schedule Builder generates a schedule corresponding to an input solution sequence. This solution is stored in a file named *tourdata* and is produced by the GENITOR code at the end of a run.

The output of the Schedule Builder code is an overlap file and a schedule file. The *overlap* file contains the start and end times of all supports which overlap into the next day or within the maximum turn around time of the end of the day. The *schedule* file contains the schedule produced for the 24-hour period. Included in this file is the support number, RTS, beginning time, ending time, service time, and turnaround time.

```
/* Schedule Builder Code - build schedule */
/* 1 Feb: print schedule to file */
/* 30 Jan 94: fix overlap by clearing the global overlaps */
/* 28 Jan 94: overlap of 20 minutes allowed at beginning */
/* 3 Jan 94: allowed supports to be scheduled past 1470 */
/* changed filled[] array from 1470 to 2100, take off condition */
#include <stdio.h>
#include <math.h>
#include <gene.h>
#include <op_edge_recomb.h>
#include <ga_global_extern.h>
int
ScheduleWindow (spt, win, filled)
int
      spt;
int
      win;
int filled[21][2100];
     /* Declare Variables */
     int
                                         gTurnAroundTime : 0 if latter <0 */
         bVis,
                           /* gBeginV_
         eVis,
                          /* ending visibility 28 Jan */
                        /* RegLength + TurnAroundTime */
         rLength,
         endSpt,
                     /* Flag for Support either scheduled or run out */
        j,
                         /* Counter for support index being scheduled */
         1,
                         /* Counter for filling in scheduled blocks */
                /*
                    spt:
                                     Index for support */
                          /* current time interval */
        minute,
                     /* length so far */
        length,
         scheduled,
                    /* binary: 1 or scheduled; 0 for not schedule */
        beginSched,
         endSched,
                    /* current start of support */
        start,
 tat,
             /* turn around time */
                     /* RTS number -- passed from global gRTS[support] */
        rts;
                       /* scheduledPtr : indicates if scheduled or not */
FILE *fp;
     /**** Initialize Variables *****/
     scheduled = 0;
     length = 0;
```

```
endSpt =0; /* initialize flag; end when it = 1*/
      /* Open file for output to schedule 1 Feb */
fp = fopen( schedule ,
                          a);
     /* Initialize Starting Point, Required Length with TAT */
    rts = gRTS[spt][win];
     tat = gTurnAroundTime[spt];
     bVis = gBeginVis[spt][win] - gTurnAroundTime[spt] + 20;
     eVis = gEndVis[spt][vin] + 20; /* 20 is max TAT - 28 Jan */
/* if ( (spt==1) || (spt==155) )
printf( tat=%d,bVis=%d\n ,tat,bVis); */
     if (bVis < 0)
£
           tat = bVis + tat;
  bVis = 0;
                   /**** reduces tat if near boundary at start ****/
}
    rLength = gReqLength[spt] + tat;
/* if ( (spt==1) || (spt==155) )
printf( rLength=%d\n ,rLength); */
    minute= bVis + 1; /* my blocks are filled to the right of position */
     start = minute;
                                /* Starting position of current attempt */
     while (endSpt != 1)
                            /* schedule until successful or end */
     £
          if (filled[rts][minute] == 0)
                                           /* if current space free */
              length += 1;
                                         /* increment length if space free */
          else
               length = 0; /* reset to zero if space is filled */
  tat = gTurnAroundTime[spt];
         rLength = gReqLength[spt] + tat;
               start = minute+1;
          }
          if (length == rLength)
               Ł
                  beginSched = minute - gReqLength[spt] - 20;
           endSched = minute - 20;
                   scheduled = 1;
                   for (l = start;l <= minute;l++)</pre>
                      filled[rts][l] = 1;
                   /* If ends in TAT zone then save to overlap */
                                    /* 1440 - max TAT; in overlap */
           if (endSched > 1420)
              if ( (gOverlapEnd[rts] = 0) || (endSched >
 gOverlapEnd[rts]) )
        { /* if rts overlap is new or this one is longer than
 last */
     gOverlapEnd[rts] = endSched - 1420;
     if (beginSched > 1420)
          gOverlapBegin[rts] = beginSched - 1420;
       gOverlapBegin[rts] = 0;
}
     else
               }
          if ( (minute== eVis) || (scheduled == 1) )
             endSpt = 1;
          minute += 1;
     } /* End Support */
      /* fprintf (stdout, scheduled = %d\n , scheduled); */
      if(scheduled == 1)
Ł
         printf( %d %d %d %d %d %d %d %
                    spt, rts, beginSched, endSched , gReqLength[spt],
```

```
gTurnAroundTime[spt] );
spt, rts, beginSched, endSched, gReqLength[spt],
spt,
gTurnAroundTime[spt] );
}
     /*else
       printf(
                    Support %d not scheduled in window %d\n , spt, win); */
    fclose(fp); /* stop printing to file 1 Feb */
    return(scheduled);
} /* end ScheduleWindow */
/***************
                        int
ScheduleSupport(spt,filled)
int spt:
int filled[21][2100];
int i, win, flag;
            /* flag = 1 if scheduled , 0 ow */
int sched;
sched = 0:
win = 1;
                  /* counter for number of windows tried */
flag = Ó;
while (flag != 1)
  £
    sched = ScheduleWindow(spt,win,filled);
                                          , sched,spt,gNumWin[spt]); */
   /* printf ( sched=%d, gWunWin[%d]=%d\n
    if ( (sched == 1) || (win == gNunWin[spt]) )
flag =1;
    win +=1;
/* if (sched==0)
 printf( *** Support %d NOT SCHEDULED!***\n ,spt); */
return(sched);
3
float
tsp_eval (order, num_supports)
int order[];
         num_supports;
int
£
    /* Declare Variables */
    int
        filled[21][2100].
                          /* time units in schedule are filled[rts][time] */
                        /* Counter for support index being scheduled */
        j,
        ž,
                      /* Index for support */
        spt,
        numScheduled.
                        /* number of supports scheduled */
         minute,
                         /* counter for current time */
        rtsNum;
                        /* max number of RTSs */
    /**** Initialize Variables *****/
   rtsNum = 20; /* must be one less than max cause of arrays 0-19 */
    /* printf( rtsNum = %d\n , rtsNum); */
   for (k = 1;k <= rtsHum;k++)
    Ł
       for (minute = 0;minute <= 2100;minute++)</pre>
           filled[k][minute] = 0;
          /* printf( filled[%d] [%d] = %d\n ,k,minute,filled[k][minute]); */
       }
    7
  /* Schedule the overlaps first, before the real supports */
```

```
for (k = 1; k \leq rtsTum; k++)
  for (minute = gOverlapBegin[k]; minute <= gOverlapEnd[k]; minute++)</pre>
filled[k][minute] = 1;
  /* Set the Overlaps back to zero; since written to in scheduling 30 Jan */
     for (k = 1; k \leq rtslum; k++)
  gOverlapBegin[k] = 0;
          gOverlapEnd[k] = 0;
     }
    numScheduled = 0;
    /* Schedule each Support in Priority Order */
printf( Spt RTS Beg End Length TAT\n\n );
    for (j = 1; j<= num_supports; j++)</pre>
                                     /* start at 0 for Tsp, 1 for sb */
    £
         spt = order[j];
                      Start SS Order[%d] = %d\n ,j, order[j]); */
            /* printf(
         numScheduled += ScheduleSupport(spt, filled);
          /* end num_supports */
   /* fprintf (stdout, num_supports = %d\n , num_supports); */
   /* The GA wants to MINimize something so do Max poss - actual */
     numScheduled);
    return(num_supports - numScheduled);
7
     /* End Eval */
                      *****************
/*********
 FUNCTION: read_data
 * DESCRIBE: read in data
 * INPUT PARAMETERS: filename of list of satellite support data
 * RETURN VALUE: none
 * CALLS: get_coord_array
         read_coords
get_2D_dist_array
         calc_distances
 *
             *****
void
read_data (coord_file, num_supports)
              coord_file[];
char
int
             num_supports;
{
  int i:
  int spt;
  int rts;
  int bVis, eVis, rLength, tat;
  FILE *fp;
  for (i=1; i<=num_supports; i++)</pre>
Ł
gWumWin[i] = 0;
                   /* set number of windows [spt] = 0 */
/* printf( gWumWin[%d]=%d\n ,i,gWumWin[i]); */
  if (fp = fopen(coord_file, r ))
{
        while (fscanf(fp, %d %d %d %d %d %d , &spt, &rts, &bVis, &eVis, &rLength, &tat ) !=EOF)
        {
             gNumWin[spt] +=1;
             gRTS[spt][gWumWin[spt]] = rts;
             gBeginVis[spt][gNumWin[spt]] = bVis;
```

```
gEndVis[spt][gNumWin[spt]] = eVis;
             gReqLength[spt] = rLength;
             gTurnAroundTime[spt] = tat;
             /* printf ( spt = %d, rts= %d, Bvis= %d, Evis = %d, RegLen = %d,
TAT = Xd\n , spt, gRTS[spt][gWunWin[spt]], gBeginVis[spt][gWunWin[spt]],
gEndVis[spt][gHumWin[spt]], gReqLength[spt], gTurnAroundTime[spt] ); */
fclose (fp);
/* for (i=1; i<= num_supports;i++)</pre>
printf( gHumWin[%d] = %d\n ,i,gHumWin[i]); */
else
fatal_error ( Cannot read input city file );
}
FUNCTION: read_overlap_data
 * DESCRIBE: read in data
 * IMPUT PARAMETERS: filename of list of satellite support overlap data
* RETURN VALUE: none
* CALLS: get_coord_array
         read_coords
         get_2D_dist_array
-
         calc_distances
******
                          void
read_overlap_data (coord_file)
char
              coord_file[];
Ł
 int i;
 int rts;
                                                         int bVis, eVis;
 FILE *fp;
 if (fp = fopen(coord_file, r ))
Ł
        while (fscanf(fp, %d %d %d , &rts, &bVis, &eVis ) !=EOF)
        Ł
             gOverlapBegin[rts] = bVis;
     gOverlapEnd[rts] = eVis;
              /* printf( g0verlapBegin[%d] = %d\t g0verlapEnd[%d] = %d\n ,
rts, gOverlapBegin[rts], rts, gOverlapEnd[rts]); */
             /* printf ( spt = %d, rts= %d, Bvis= %d, Evis = %d, ReqLen = %d,
 TAT = %d\n , spt, gRTS[spt][gHumWin[spt]], gBeginVis[spt][gHumWin[spt]],
gEndVis[spt][gHumWin[spt]], gReqLength[spt], gTurnAroundTime[spt] ); */
TAT = %d n
     /* printf( -----gNumWin[%d] = %d\n , spt,gNumWin[spt]); */
 fclose (fp);
  /* for (i=1; i<= num_supports;i++)</pre>
printf( gHumWin[%d] = %d\n ,i,gHumWin[i]); */
 else
fatal_error ( Cannot read input city file );
}
/*****
                        ****************
FUNCTION: print_overlap_data
 * DESCRIBE: print out data
```

```
* IMPUT PARAMETERS: filename of list of satellite support overlap data
 ***
RETURN VALUE: none
 void
print_overlap_data (coord_file)
char
               coord_file[];
{
    int i;
    int rts;
    int bVis, eVis;
    FILE *fp;
 if (fp = fopen(coord_file, v ))
£
     for (rts = 1; rts <= 19; rts++)
{</pre>
                 bVis = gOverlapBegin[rts];
          eVis = gOverlapEnd[rts];
fprintf(fp, %d\t%d\n , rts, bVis, eVis );
             7
     fclose (fp);
}
else
fatal_error ( Cannot read or write overlap file );
}
```

Appendix D. Satellite Range Scheduling Data Processing

The processing of satellite support data begins with raw ASTRO data which is processed by Pascal programs to produce satellite request windows. This data is filtered by a C program to prepare the data for the genetic algorithm.

ASTRO Data

ASTRO (Automated Scheduling Tools for Range Operations) is a computer system and database to aid the range schedulers. The report information for the seven days of test data is in a file named FINLDATA.DFT. This file is the raw input which determines support requests and request visibilities.

Requests and Visibility Windows

Pascal programs written by Gooley and Schalck are used to process the ASTRO data. The programs extract the satellite support requirements and request visibilities. Five programs are used for this algorithm: one to process the low-altitude satellite (lowflyer) support information, and four to process the medium and high-altitude satellite support (high-flyer) information. These programs are run using TurboPascal on a IBM PC compatible computer.

The final data format is:

- 1. Support Number
- 2. RTS the satellite is visible to
- 3. Beginning visibility time of window (in minutes)
- 4. Ending visibility time of window
- 5. Length of support (in minutes)
- 6. TAT (turn around time) required by an RTS
- 7. IRON/Revolution: identifies satellite and pass

LREQ.PAS. This Pascal program reads in the information on low-flyer supports and saves the information to a file called REQLF.DAT containing the low altitude satellite support requests for a day.

HREQ.PAS. This program reads in the information on high-altitude requests for a day and saves the information to two files: 1) REQHF.DAT: the high-altitude requests for a day, and 2) D1V.DAT: the visibilities for the high altitude requests for a day.

TOL.PAS. This program takes tolerance data presented in different formats for each request from HREQ.PAS and standardizes the tolerance window data.

CROSS2.PAS. This program cross checks requests and visibilities. It cross references the visibility file created in HREQ.PAS and the output from TOL.PAS to determine all the RTSs that can satisfy each medium or high altitude support request from TOL.PAS.

RTS.PAS. This Pascal program ensures all RTS sides are included once and only once for each support request-RTS visibility combination.

Prepare Time Window Data for GENITOR

This program reads in the separate data files for the low-flyers and high-flyers and combines them into one file for use in the GENITOR genetic algorithm. It also changes the RTS name to a number for reference in array structures.

```
/* 5 Jan 94 Preprocess Data modications */
/* now strips out all windows for RTS # 9, 16, 19 to match Spike */
/* -- reads in request windows of LF/HF and prints out concate file */
/*
        with RTS names changed to numbers
/* 1) Open output file "Dayout.dat"
                                    */
/* 2) Ask for LF file name ( ex. "LFDay1g.dat") */
/* 3) Read in data, convert name to number and drop IRON */
/* 4) Ask for HF file name, read in data -- start spt # at LF +1 */
                   *********
                                                                *******
Program: Process.
 * DESCRIBE: read in data
 * IMPUT PARAMETERS: filename of list of satellite support data for LF and HF
* RETURN VALUE: none, but print a support list to Dayout.dat
 * CALLS: read_data
```

```
#include <stdio.h>
int
equal_strings (s1, s2)
char s1[], s2[];
int i = 0, answer;
while ( s1[i] == s2[i] && s1[i] != '\0' && s2[i] != '\0' ) ++i;
if ( s1[i] == '\0' && s2[i] == '\0' )
answer = 1; /* strings equal */
else
answer = 0; /* not equal
                              */
return (answer);
}
int
name_to_Wum (name)
char name[];
int rts;
        (equal_strings (name, "POGO-A") )
iŕ
rts = 1;
else if (equal_strings (name, "POGO-B") )
rts = 2;
else if (equal_strings (name, "POGO-C") )
rts = 3;
else if (equal_strings (name, "HULA-A") )
rts = 4;
else if (equal_strings (name, "HULA-B") )
rts = 5;
else if (equal_strings (name, "COOK-A") )
rts = 6;
else if (equal_strings (name, "COOK-B") )
rts = 7;
else if (equal_strings (name, "INDI-A") )
rts = 8;
else if (equal_strings (name, "INDZ-B") )
rts = 0;
else if (equal_strings (name, "BOSS-A") )
rts = 10;
else if (equal_strings (name, "BOSS-B") )
rts = 11:
else if (equal_strings (name, "LIOM-A") )
rts = 12;
else if (equal_strings (name, "LION-B") )
rts = 13;
else if (equal_strings (name, "GUAM-A") )
rts = 14;
else if (equal_strings (name, "GUAM-B") )
rts = 15:
else if (equal_strings (name, "PIKE-A") )
rts = 16:
else if (equal_strings (name, "PIKE-B") )
rts = 0;
else if (equal_strings (name, "REEF-A") )
rts = 18;
else if (equal_strings (name, "REEF-B") )
rts = 0;
else
printf("UNKNOWN\n");
return(rts);
}
read_data (out_file, coord_file, start, endPtr)
FILE *out_file;
```

```
char
            coord_file[];
int start;
int *endPtr;
int i;
int spt;
  char rtsWame[8];
int bVis, eVis, rLength,tat;
   char iron[16];
   int totNumSpt;
  int rts;
FILE *input_file;
   input_file = fopen(coord_file, "r");
  while (fscanf(input_file, "%d %s %d %d %d %d %s", &spt, &rtsWame, &bVis,
 &eVis, &rLength, &tat, &iron ) !=EOF)
 spt = spt + start;
 totNumSpt = spt;
 rts = name_to_Num(rtsName);
if ( rts != 0 )
            rLength, tat);
    }
   fclose (input_file);
   *endPtr = totWumSpt;
}
/****
          MAIN PROGRAM *****/
main()
{ /* start MAIN */
/**** Initialize and declare Variables *****/
char lfdata[12], hfdata[12];
int WumLFSpt, WumHFSpt;
FILE *output_file;
output_file = fopen("datafile", "w");
/**** Input LF Filename *****/
printf("What is the LF filename?");
scanf ("%s", lfdata);
/**** Read in LF Data and print to Dayout.dat *****/
read_data (output_file, lfdata, 0, &WumLFSpt);
/**** Input HF filename *****/
printf("What is the HF filename?");
scanf ("%s", hfdata);
/**** Read in HF Data and print to Dayout.dat *****/
read_data (output_file, hfdata,MumLFSpt, &MumHFSpt);
fclose (output_file);
printf("WumLFSpt = %d, EndHFSpt = %d\n", WumLFSpt, WumHFSpt);
} /* end Main */
```

Appendix E. Schedules for Week One Data

Schedules produced for each of seven days of data follow. The scheduled supports are sorted by RTS, and by time. The columns are: Support number, RTS, Beginning Time, Ending Time, Support Length, and Support turn around time.

Day 1

Day 1
2 POGO-A 26 42 16 20 167 POGO-A 65 80 15 15 170 POGO-A 95 105 10 15 10 POGO-A 138 152 14 20
167 POGO-A 65 80 15 15 170 POGO-A 95 105 10 15
10 POGO-A 138 152 14 20 180 POGO-A 180 185 5 15 22 POGO-A 234 248 14 20 27 POGO-A 268 279 11 20 10 POGO-A 168 10 15
193_P0G0-A_294_304_10_15
100 F000-A 100 108 10 22 P000-A 234 248 14 20 27 P000-A 268 279 11 20 193 P000-A 294 304 10 15 31 P000-A 331 345 14 20 39 P000-A 367 379 12 20 204 P000-A 367 379 12 20
27 PUGU-A 206 279 11 20 193 POGU-A 294 304 10 15 31 POGU-A 331 345 14 20 39 POGU-A 367 379 12 20 204 POGU-A 367 379 12 20 204 POGU-A 395 410 15 15 205 POGU-A 425 435 10 15 47 POGU-A 426 496 14 20
47 POGO-A 466 480 14 20 51 POGO-A 527 543 16 20
51 POGO-A 527 543 16 20 57 POGO-A 567 583 16 20
22 POGO-A 234 248 14 20 27 POGO-A 268 279 11 20 193 POGO-A 367 379 12 20 204 POGO-A 367 379 12 20 204 POGO-A 395 410 15 15 205 POGO-A 425 435 10 15 47 POGO-A 466 480 14 20 51 POGO-A 667 543 16 20 228 POGO-A 667 543 16 20 235 POGO-A 690 700 10 15 240 POGO-A 690 700 10 15 240 POGO-A 725 735 10 15 76 POGO-A 756 772 16 20 248 POGO-A 800 835 35 15 259 POGO-A 800 835 35 15 269 POGO-A 1020 1035 15 15 99 POGO-A 1020 1035 15 15 109 POGO-A 1089 1099 10 15 116 POGO-A 1221 1235 14 20 127 POGO-A 1359 1375 16 20 276 POGO-A 1318 1333 15 20 124 POGO-A 1318 1333 15 20 145 POGO-A 1359 1375 16 20 322 POGO-A 1415 1450 35 15 158 POGO-A 136 120 15 4 POGO-B 51 64 13 20 172 POGO-B 105 120 15 158
235 POGO-A 690 700 10 15 240 POGO-A 725 735 10 15
76 PDG0-A 756 772 16 20
248 POGO-A 800 835 35 15 259 POGO-A 870 890 20 15
94 POGO-A 939 950 11 20 99 POGO-A 971 987 16 20
276 POGO-A 1020 1035 15 15 109 POGO-A 1062 1074 12 20
280 POGO-A 1089 1099 10 15 116 POGO-A 1123 1136 13 20
116 POGO-A 1123 1136 13 20 124 POGO-A 1171 1188 17 20
259 F0G0-A 970 390 950 11 20 99 P0G0-A 971 987 16 20 276 P0G0-A 1020 1035 15 15 109 P0G0-A 1062 1074 12 20 280 P0G0-A 1089 1099 10 15 116 P0G0-A 1123 1136 13 20 124 P0G0-A 1171 1188 17 20 127 P0G0-A 1260 1274 14 20 134 P0G0-A 1318 1333 15 20 135 P0G0-A 1359 1375 16 20 322 P0G0-A 1415 1450 35 15 158 P0G0-B 5 15 10 15
134 POGO-A 1260 1274 14 20 139 POGO-A 1318 1333 15 20 145 POGO-A 1359 1375 16 20
145 POGO-A 1359 1375 16 20 322 POGO-A 1415 1450 35 15
158 POGO-B 5 15 10 15 4 POGO-B 51 64 13 20 172 POGO-B 105 120 15 15
172 POGO-B 105 120 15 15 15 POGO-B 169 179 10 20
15 POGO-B 169 179 10 20 21 POGO-B 232 248 16 20 200 POGO-B 350 360 10 15
20 POGO-B 232 248 16 20 200 POGO-B 350 360 10 15 203 POGO-B 375 390 15 15
42 PUGU-B 420 443 17 20 200 POGO-R 475 520 45 15
56 P0G0-B 559 575 16 20 62 P0G0-B 625 639 14 20 67 P0G0-B 660 674 14 20 236 P0G0-B 690 725 35 15 77 P0G0-B 757 771 14 20 236 P0G0-B 757 771 14 20
62 POGO-B 625 639 14 20 67 POGO-B 660 674 14 20 236 POGO-B 690 725 35 15
77 POGO-B 757 771 14 20 83 POGO-B 834 848 14 20
k3 POGO-B 834 848 14 20 252 POGO-B 863 878 15 15 263 POGO-B 915 960 45 15 267 POGO-B 975 990 15 15 274 POGO-B 1010 1020 10 15 107 POGO-B 1047 1061 14 20 284 POGO-B 1080 1095 15 15 289 POGO-B 1110 1125 15 15 119 POGO-B 1145 1158 13 20
263 POGO-B 915 960 45 15 267 POGO-B 975 990 15 15
83 P0G0-B 634 648 14 20 252 P0G0-B 863 878 15 15 263 P0G0-B 915 960 45 15 267 P0G0-B 975 990 15 15 274 P0G0-B 1010 1020 10 15 107 P0G0-B 1047 1061 14 20 284 P0G0-B 1080 1095 15 15 366 P0G0-B 1110 1125 15
284 POGO-B 1080 1095 15 15
289 POGO-B 1110 1125 15 15 119 POGO-B 1145 1158 13 20
289 POGO-B 1110 1125 15 15 119 POGO-B 1145 1158 13 20 296 POGO-B 1145 1158 13 10 293 POGO-B 1198 1208 10 15 131 POGO-B 1247 1262 15 20
289 POGO-B 1110 1125 15 15 119 POGO-B 1145 1158 13 20 296 POGO-B 1145 1158 13 20 293 POGO-B 1198 1208 10 15 131 POGO-B 1247 1262 15 20 308 POGO-B 1290 1330 40 15
14 DOCO_D 1940 1975 19 00
168 P0G0-C 150 163 15 15 188 P0G0-C 150 163 15 20 182_P0G0-C 205_220 15 15
13 POGO-C 100 103 13 20 182 POGO-C 205 220 15 15 26 POGO-C 249 264 15 20 30 POGO-C 326 342 16 20 37 POGO-C 364 376 12 20 45 POGO-C 449 466 17 20
26 P0G0-C 249 264 15 20 30 P0G0-C 326 342 16 20 37 P0G0-C 364 376 12 20
211 POGO-C 505 515 10 15 54 POGO-C 551 567 16 20 224 POGO-C 582 622 40 15
224 POGO-C 582 622 40 15 66 POGO-C 653 669 16 20 232 POGO-C 684 699 15 15
54 PUGU-C 551 557 15 20 224 POGU-C 582 622 40 15 66 PUGU-C 653 669 16 20 232 PUGU-C 684 699 15 15 75 PUGU-C 739 753 14 20
88 PUGU-C 854 868 14 20
262 PUGU-C 910 955 45 15 266 PUGU-C 970 980 10 15
103 P060-C 1024 1036 12 20 111 P060-C 1071 1087 16 20
290 POGO-C 1115 1130 15 15
299 POGO-C 1200 1245 45 15 305 POGO-C 1265 1290 25 15
305 POGO-C 1265 1290 25 15

142 POGO-C 1335 1351 16 20
154 HULA-A 0 35 35 15 161 HULA-A 50 70 20 15
169 HULA-A 90 105 15 15 12 HULA-A 139 153 14 20 175 HULA-A 168 188 20 15
184 HULA-A 210 225 15 15 185 HULA-A 240 245 5 15 32 HULA-A 335 351 16 20
198 HULA-A 366 391 25 15 40 HULA-A 413 430 17 20
40 HULA-A 413 430 17 20 50 HULA-A 515 527 12 20 55 HULA-A 551 561 10 20 225 HULA-A 590 600 10 15
226 HULA-A 615 620 5 15 230 HULA-A 660 675 15 15
230 HULA-A 660 675 15 15 74 HULA-A 736 752 16 20 244 HULA-A 767 772 5 15 80 HULA-A 807 820 13 20
249 HULA-A 835 850 15 15 255 HULA-A 865 880 15 15 89 HULA-A 904 917 13 20
264 HULA-A 932 947 15 15 98 HULA-A 970 981 11 20
264 HULA-A 932 947 15 15 98 HULA-A 970 981 11 20 270 HULA-A 996 1016 20 15 277 HULA-A 1035 1040 5 15 112 HULA-A 1079 1095 16 20 288 HULA-A 1110 1120 10 15 266 HULA-A 1110 1120 10 15
112 HULA-A 1079 1095 16 20 288 HULA-A 1110 1120 10 15 295 HULA-A 1155 1160 5 15 298 HULA-A 1155 1160 5 15
295 HULA-A 1155 1160 5 15 298 HULA-A 1190 1205 15 15 304 HULA-A 1260 1285 25 15 312 HULA-A 1325 1335 10 15
315 HULA-A 1355 1368 13 15 160 HULA-B 30 75 45 15 24 HULA-B 238 254 16 20
24 HULA-B 238 254 16 20 33 HULA-B 335 351 16 20 199 HULA-B 366 371 5 15
212 HULA-B 510 530 20 15 221 HULA-B 560 575 15 15 73 HULA-B 736 752 16 20 239 HULA-B 767 772 5 15
239 HULA-B 767 772 5 15 81 HULA-B 807 820 13 20 254 HULA-B 855 900 45 15
97 HULA-B 970 981 11 20 102 HULA-B 1020 1036 16 20 272 HULA-B 1051 1056 5 15
272 HULA-B 1051 1056 5 15 113 HULA-B 1079 1095 16 20 123 HULA-B 1169 1181 12 20 128 HULA-B 1222 1234 12 20 133 HULA-B 1202 1234 12 20 133 HULA-B 1207 1268 11 20 309 HULA-B 1290 1305 15_15_
128 HULA-B 1222 1234 12 20 133 HULA-B 1257 1268 11 20 309 HULA-B 1290 1305 15 15
310 HULA-B 1320 1485 165 15 156 COOK-A 0 5 5 15 5 COOK-A 0 5 5 15
142 POGO-C 1335 1351 16 20 154 HULA-A 0 35 35 15 161 HULA-A 90 105 15 15 12 HULA-A 139 153 14 20 175 HULA-A 139 153 14 20 175 HULA-A 168 188 20 15 184 HULA-A 210 225 15 155 HULA-A 335 351 16 20 198 HULA-A 366 391 25 15 40 HULA-A 366 391 25 15 226 HULA-A 615 620 5 15 230 HULA-A 615 620 5 15 230 HULA-A 615 620 5 15 15 HULA-A 767 772 5 15 80 HULA-A 767 772 5 15 80 HULA-A 767 772 5 15 12 HULA-A 767 772 5 15 12 HULA-A 994 917 13 20 244 HULA-A 994 917 13 20 249 HULA-A 996 1016 20 15 277 HULA-A 996 1016 20 15 277 HULA-A 1035 1040 5 15 112 HULA-A 1079 1095 16 20 286 HULA-A 1155 1160 5 15 304 HULA-A 1325 1368 13 15 160 HULA-A 1325 1368 13 15 160 HULA-A 1325 1368 13 15 160 HULA-A 970 981 11 20 270 HULA-A 1325 1368 13 15 160 HULA-A 1155 1160 5 15 312 HULA-A 1325 1368 13 15 160 HULA-A 1325 1368 13 15 160 HULA-B 30 75 45 15 315 HULA-B 135 146 5 15 315 HULA-B 135 146 5 15 315 HULA-A 1325 1368 13 15 160 HULA-B 1079 1095 16 20 288 HULA-A 1150 1206 5 15 315 HULA-A 1260 1285 25 15 315 HULA-A 1260 1285 25 15 315 HULA-A 1260 1285 25 15 315 HULA-B 1355 1368 13 15 160 HULA-B 1355 1368 13 15 160 HULA-B 135 146 5 15 214 HULA-B 1355 146 20 199 HULA-B 1355 146 20 199 HULA-B 1355 146 20 199 HULA-B 1355 146 15 211 HULA-B 1355 10 15 213 HULA-B 1355 10 15 315 HULA-B 1451 1056 5 15 314 HULA-B 1451 1056 5 15 315 HULA-B 1451 1056 5 15 314 HULA-B 169 1181 12 00 102 HULA-B 169 1181 12 00 102 HULA-B 169 1181 12 00 103 HULA-B 169 1181 12 00 104 HULA-B 169 1181 12 00 105 HULA-B 169 1181 12 00 107 HULA-B 169 1181 12 00 107 HULA-B 169 1181 12 00 108 HULA-B 169 1181 12 00 109 HULA-B 169 1181 12 00 102 HULA-B 169 1181 12 00 103 HULA-B 169 1181 12 00 104 HULA-B 169 1181 12 00 105 HULA-B 169 1181 12 00 107 HULA-B 1290 1306 16 15 131 HULA-B 169 1181 12 00 123 HULA-B 169 1181 12 00 124 HULA-B 169 1181 12 00 1254 HULA-B 169 1181 12 00 126 HULA-B 169 1181 12 00 127 HULA-B 169 1185 15 150 140 HULA-B 169 155 15 150 140 HULA-B 160 15 15 151 150 140 HULA-B 160 15 15 152 19 HULA-B 160 1
194 COOK-A 300 345 45 15 44 COOK-A 441 452 11 20 216 COOK-A 525 535 10 15
219 COOK-A 550 575 25 15 231 COOK-A 660 705 45 15 242 COOK-A 735 765 30 15
88 CODK-A 874 890 16 20
261 COOK-A 905 920 15 15 100 COOK-A 975 989 14 20 282 COOK-A 1080 1095 15 15 283 COOK-A 1110 1130 20 15
138 CUUR-A 1293 1308 15 20 316 COOK-A 1365 1370 5 15
6 COOK-B 54 68 14 20 26 COOK-B 54 68 14 20
214 COOK-B 515 525 10 15 241 COOK-B 730 925 195 15 297 COOK-B 1185 1192 7 15
136 COOK-B 1274 1290 16 20 311 COOK-B 1320 1330 10 15
104 1901-8 04 04 10 10
171 INDI-A 90 110 20 15 17 INDI-A 179 195 16 20 183 INDI-A 210 250 40 15

Day 2

	.5
157 POGO-A 150 170 20 1	5
162 POGO-A 200 215 15 1	5
166 POGO-A 230 280 50 1	15
173 POGO-X 295 305 10 1	ĪŠ
180 POGO-A 345 390 45 1	15
31 POGO-A 419 435 16 20	
36 POGO-A 460 475 15 20	ί.
36 POGO-A 460 475 15 20 41 POGO-A 512 529 17 20	ί.
45 POGO-A 554 570 16 20	ί.
52 POGO-A 607 619 12 20	<
211 POGO-A 645 660 15 1	0
61 PDG0-A 690 705 15 20	2
222 PUGU-A 720 735 15 1	ιÞ
222 POGO-A 720 735 15 227 POGO-A 750 780 30 1	15
73 POGU-A 808 823 15 20)
78 POGO-A 858 868 10 20	
85 POGO-A 907 923 16 20)
	15

145 POGO-A 45 90 45 15

94 POGO-A 1015 1027 12 20 262 POGO-A 1065 1075 10 15 265 POGO-A 1090 1105 15 15 103 POGO-A 1132 1147 15 20 272 POGO-A 1130 1190 10 15 113 POGO-A 1232 1246 14 20 280 POGO-A 1331 1346 15 20 298 POGO-A 1331 1346 15 20 298 POGO-A 1330 1395 15 142 POGO-B 135 150 15 15 149 POGO-B 65 100 35 15 160 POGO-B 135 150 15 15 170 POGO-B 270 285 15 15 170 POGO-B 270 285 15 15 177 POGO-B 300 345 15 181 POGO-B 300 375 15 15 181 POGO-B 385 410 15 15 181 POGO-B 385 410 15 15 187 POGO-B 395 410 15 15 181 POGO-B 395 410 15 15 177 POGO-B 395 410 15 15 186 POGO-B 482 497 15 15 197 POGO-B 737 801 14 20 196 POGO-B 737 801 14 20 204 POGO-B 736 746 10 15 71 POGO-B 737 801 14 20 74 POGO-B 133 1145 12 20 207 POGO-B 133 1145 12 20 21 POGO-B 133 1145 12 20 21 POGO-B 136 133 1145 12 20 221 POGO-B 136 133 1145 12 20 232 POGO-B 136 132 1378 16 20 242 POGO-B 133 1145 12 20 271 POGO-B 1304 1320 16 20 232 POGO-B 1304 1320 16 20 239 POGO-C 0 10 10 15 184 POGO-B 1233 1247 14 20 139 POGO-C 210 240 30 15 138 POGO-C 100 10 15 184 POGO-C 210 240 30 15 175 POGO-C 210 240 30 15 184 POGO-C 135 355 10 15 189 POGO-C 410 455 45 15 189 POGO-C 410 455 45 15 189 POGO-C 410 456 45 15 189 POGO-C 410 456 45 15 189 POGO-C 113 1143 122 0 235 POGO-C 158 607 10 15 189 POGO-C 1131 1143 122 0 235 POGO-C 158 1075 17 20 102 POGO-C 840 845 5 15 138 POGO-C 1351 1366 15 15 243 POGO-C 1353 1366 13 10 257 POGO-C 1351 1366 15 15 243 POGO-C 1351 1366 15 15 244 POGO-C 1351 1366 15 15 245 POGO-C 1355 155 15 245
$ \begin{array}{c} 33 \\ 196 \\$
217 FUGD-C 758 769 11 20 235 POGD-C 758 769 11 20 235 POGD-C 860 875 15 15 243 POGD-C 905 945 40 15 257 POGD-C 1005 1015 10 15 98 POGD-C 1058 1075 17 20 102 POGD-C 1131 1143 12 20 268 POGD-C 1151 1143 12 20 268 POGD-C 1265 1221 16 20 117 POGD-C 1260 1276 16 20 117 POGD-C 1260 1276 16 20 123 POGD-C 1323 1336 13 20 293 POGD-C 1351 1366 15 15 134 POGD-C 1420 1435 15 20 4 HULA-A 40 54 14 20 144 HULA-A 69 89 20 15 154 HULA-A 120 135 15 15 151 HULA-A 301 30 10 15 154 HULA-A 383 398 15 20 167 HULA-A 383 398 15 20 183 HULA-A 435 450 15 15 199 HULA-A 540 530 20 15 199 HULA-A 510 530 20 15 202 HULA-A 510 530 20 15 202 HULA-A 706 722 16 20 215 HULA-A 706 722 16 20 224 HULA-A 706 722 16 20 225 HULA-A 750 755 5 15 229 HULA-A 750 755 5 15 229 HULA-A 815 820 5 15 229 HULA-A 815 820 5 15 238 HULA-A 865 885 20 15 334 HULA-A 865 885 20 15 34 HULA-A 865 885 20 15 35 H

.

248 HULA-A 945 950 249 HULA-A 965 975 254 HULA-A 965 975 254 HULA-A 1050 106 100 HULA-A 1155 11 275 HULA-A 1155 11 275 HULA-A 1200 12 288 HULA-A 1200 12 288 HULA-A 1305 13 290 HULA-A 1350 13 290 HULA-A 1350 14 302 HULA-A 1350 14 302 HULA-B 160 165 515 HULA-B 160 165 515 HULA-B 160 165 515 HULA-B 482 497 209 HULA-B 482 497 209 HULA-B 482 497 209 HULA-B 1005 102 267 HULA-B 1005 102 267 HULA-B 1100 11 273 HULA-B 1185 12 285 HULA-B 185 15 219 HULA-B 1350 13 140 COOK-A 282 297 24 COOK-A 282 297 24 COOK-A 334 351 187 COOK-A 470 498 220 COOK-A 795 830 230 COOK-A 845 860	5 15 10 15 17 20 15 30 15 17 16 30 15 15 15 15 15 15 15 15 15 15 15 15 15 15 15 200 155 15 155 15
100 HULA-A 1101 11 209 HULA-A 1105 11 276 HULA-A 1200 12 288 HULA-A 1200 12 288 HULA-A 1300 14 302 HULA-A 1300 14 302 HULA-A 1300 14 302 HULA-A 1300 14 143 HULA-B 60 65 5 219 HULA-B 160 165 211 HULA-B 307 322 299 HULA-B 482 497 209 HULA-B 482 497 209 HULA-B 730 95 223 HULA-B 730 95 2245 HULA-B 1100 11 273 HULA-B 1100 11 273 HULA-B 1100 11 273 HULA-B 1185 12 285 HULA-B 1185 12 292 COOK-A 282 297 24 COOK-A 288 287 200 COOK-A 16 31 1 16 COOK-A 288 297 24 COOK-A 16 31 1 187 COOK-A 288 297 24 COOK-A 16 31 1 185 COOK-A 1020 12 236 COOK-A 795 830 236 COOK-A 1020 12 115 COOK-A 1290 13 130 COOK-A 1364 13 140 COOK-A 1290 13 130 COOK-A 1290 13 130 COOK-A 1290 13 144 COOK-B 915 960 253 COOK-A 1290 13 130 COOK-B 915 960 253 COOK-B 913 960 726 254 COOK-B 1346 13 147 COOK-B 1346 13 147 COOK-B 1346 13 147 COOK-B 1305 13 140 COOK-B 1305 13 140 COOK-B 209 226 156 COOK-B 1346 13 147 COOK-B 1305 13 140 COOK-B 1305 13 140 COOK-B 915 960 213 INDI-A 265 283 176 INDI-A 165 199 156 INDI-A 165 199 156 INDI-A 165 199 156 INDI-A 370 393 192 INDI-A 370 393 192 INDI-A 370 393 192 INDI-A 1050 1 259 INDI-A 1305 13 192 INDI-A 1305 13 193 INDI-A 1305 13 194 INDI-A 1305 13 195 INDI-A 140 144 159 INDI-A 1305 13 195 INDI-A 140 144 195 INDI-A 1305 13 195 INDI-A 1405 11 266 INDI-A 1305 13 195 INDI-A 1405 11 266 INDI-A 1405 11 266 INDI-A 1405 11 267 INDI-A 1405 11 268 INDI-A 1305	$\begin{array}{c} 25 & 15 \\ 25 & 15 \\ 15 & 15 \\ 15 & 15 \\ 15 & 15 \\ 15 & 15 \\ 16 & 10 \\ 16 & 10 \\ 10 & 16 \\ 10 & 16 \\ 10 & 16 \\ 10 & 16 \\ 10 \\ 10 & 16 \\ 10 \\ 10 \\ 10 \\ 10 \\ 10 \\ 10 \\ 10 \\$

165 BOSS-A 225 230 5 15 174 BOSS-A 270 290 20 15 22 BOSS-A 410 426 16 20 34 BOSS-A 452 466 14 20 198 BOSS-A 560 570 10 15 49 BOSS-A 664 621 17 20 55 BOSS-A 671 691 20 15 214 BOSS-A 706 716 10 15 68 BOSS-A 742 758 16 20 225 BOSS-A 940 951 11 20 250 BOSS-A 940 951 11 20 250 BOSS-A 966 981 15 15 96 BOSS-A 1165 1177 12 20 274 BOSS-A 1640 1056 16 20 107 BOSS-A 1640 1056 16 20 107 BOSS-A 1640 1056 15 282 BOSS-A 1640 1056 15 285 BOSS-B 1040 1056 15 285 BOSS-B 1040 1056 15 135 BOSS-B 602 616 14 20 208 BOSS-B 603 60 300 15 38 BOSS-B 1070 1090 20 15 281 BOSS-B 1270 1295 25 15 289 BOSS-B 1320 1485 165 15 244 BOSS-B 1320 1485 165 15 245 BOSS-B 1320 1485 165 15 244 BOSS-B 1320 1485 165 15 245 BOSS-B 1320 1485 165 15 249 BOSS-B 1320 1485 165 15 251 LOW-A 136 155 20 15 13 LIOW-A 257 271 14 20 25 LIOW-A 336 351 15 20 25 LIOW-A 337 453 16 20 25 LIOW-A 337 749 12 20 76 LIOW-A 877 852 15 20 81 LIOW-A 397 1985 14 20 25 LIOW-A 397 1985 14 20 25 LIOW-A 1209 1219 10 15 287 LIOW-A 397 1985 14 20 257 LIOW-A 1209 1219 10 15 287 LIOW-A 1209 1219 10 15 299 LIOW-B 45 480 435 15 201 LIOW-B 655 670 15 15 201 LIOW-B 655 670 15 15 201 LIOW-B 655 670 15 15 201 LIOW-B 1223 1239 16 20 299 LIOW-B 965 980 15 20 99 LIOW-B 965 980 15 20 99 LIOW-B 965 980 15 20 99 LIOW-B 1063 1075 12 20 20 111 LIOW-B 1263 1075
22 BOSS-A 310 324 14 20 30 BOSS-A 410 426 16 20 34 BOSS-A 452 466 14 20
198 BOSS-A 510 520 10 15 205 BOSS-A 560 570 10 15 49 BOSS-A 604 621 17 20
55 BOSS-A 641 656 15 20 206 BOSS-A 671 691 20 15 214 BOSS-A 706 716 10 15
68 BOSS-A 742 758 16 20 225 BOSS-A 773 793 20 15 242 BOSS-A 900 10 10 15
87 BOSS-A 940 951 11 20 250 BOSS-A 966 981 15 15
107 BOSS-A 1165 1177 12 20 274 BOSS-A 1195 1200 5 15
282 BUSS-A 12/5 1290 16 16 297 BOSS-A 1380 1400 20 15 135 BOSS-A 1422 1438 16 20
148 BUSS-B 60 360 300 15 38 BUSS-B 505 518 13 20 48 BUSS-B 602 616 14 20
208 BUSS-B 631 636 5 15 216 BUSS-B 675 680 5 15 244 BUSS-B 915 930 15 15
260 BOSS-B 1040 1050 10 15 263 BOSS-B 1070 1090 20 15 281 BOSS-B 1270 1295 25 15
289 BOSS-B 1320 1485 165 15 2 LION-A 11 24 13 20 153 LION-A 110 120 10 15
150 LION-A 135 155 20 15 13 LION-A 209 222 13 20 18 LION-A 257 271 14 20
25 LION-A 336 351 15 20 28 LION-A 404 414 10 20 32 LION-A 437 453 16 20
42 LION-A 540 551 11 20 54 LION-A 635 646 11 20 62 LION-A 693 706 13 20
67 LION-A 737 749 12 20 76 LION-A 837 852 15 20 81 LION-A 875 888 13 20
90 LION-A 971 985 14 20 255 LION-A 1000 1020 20 15 264 LION-A 1080 1100 20 15
101 LION-A 1122 1137 15 20 276 LION-A 1209 1219 10 15 287 LION-A 1295 1305 10 15
131 LION-A 1369 1384 15 20 136 LION-A 1422 1433 11 20 300 LION-A 1448 1458 10 15
146 LION-B 45 480 435 15 201 LION-B 525 550 25 15 212 LION-B 655 670 15 15
66 LION-B 735 751 16 20 80 LION-B 868 880 12 20 89 LION-B 965 980 15 20
99 LION-B 1063 1075 12 20 111 LION-B 1223 1239 16 20 284 LION-B 1280 1290 10 15 121 LION-B 1314 1330 16 20
121 LIUN-B 1314 1330 16 20
294 L108 B 143 156 13 20 9 GUAN-A 143 156 13 20 17 GUAN-A 243 259 16 20 26 GUAN-A 373 387 14 20 39 GUAN-A 507 521 14 20 51 GUAN-A 607 623 16 20
1/ GUAN-A 232 209 10 20 26 GUAN-A 373 387 14 20 39 GUAN-A 507 521 14 20 51 GUAN-A 607 623 16 20 59 GUAM-A 669 684 15 20 70 GUAM-A 766 781 15 20 233 GUAM-A 850 850 20 15
59 GUAN-A 669 684 15 20 70 GUAN-A 766 781 15 20 233 GUAN-A 830 850 20 15 83 GUAN-A 905 921 16 20 92 GUAN-A 1007 1021 14 20 108 GUAN-A 1205 1209 9 20 116 GUAN-A 1254 1270 16 20
108 GUAN-A 1200 1209 9 20 116 GUAN-A 1254 1270 16 20 127 GUAN-A 1357 1370 13 20 291 GUAN-A 1385 1405 20 15
172 GUAM-B 265 745 480 15 256 GUAM-B 1000 1660 660 15
7 PIKE-A 110 126 16 20 11 PIKE-A 167 182 15 20 40 PIKE-A 509 525 16 20
40 PIKE-A 509 525 16 20 44 PIKE-A 550 562 12 20

• •

.

Day 3

113 POGO-B 1146 1162 16 20 121 POGO-B 1218 1232 14 20 290 POGO-B 1247 1257 10 15 128 POGO-B 1247 1257 10 15 128 POGO-B 1350 1366 16 20 151 POGO-C 30 45 15 15 7 POGO-C 90 103 13 20 163 POGO-C 120 135 15 15 17 POGO-C 189 203 14 20 186 POGO-C 285 295 10 15 196 POGO-C 360 375 15 15 199 POGO-C 300 405 15 15 205 POGO-C 420 435 15 15 47 POGO-C 501 517 16 20
113 POGO-B 1146 1162 16 20 121 POGO-B 1218 1232 14 20 129 POGO-B 1247 1257 10 15 128 POGO-B 1289 1303 14 20 151 POGO-C 30 45 15 15 7 POGO-C 120 135 15 15 163 POGO-C 285 295 10 15 196 POGO-C 285 295 10 15 196 POGO-C 360 375 15 15 205 POGO-C 599 615 16 20 216 POGO-C 599 615 16 20 216 POGO-C 599 615 16 20 248 POGO-C 797 813 16 20 248 POGO-C 797 813 16 20 248 POGO-C 797 813 16 20 248 POGO-C 100 721 11 20 78 POGO-C 975 990 15 15 108 POGO-C 1101 1117 16 20 279 POGO-C 1301 1374 13 20 148 HULA-A 30 45 15 15 158 HULA-A 30 45 15 15 158 HULA-A 30 45 15 15 168 HULA-A 30 45 15 15 178 HULA-A 30 45 15 178 HULA-A 30 45 15 178 HULA-A 30 45 15 178 HULA-A 30 557 10 15 185 HULA-A 300 50 10 15 185 HULA-A 300 50 10 15 178 HULA-A 300 50 10 15 194 HULA-A 398 408 10 20 210 HULA-A 398 408 10 20 231 HULA-A 398 408 10 20 231 HULA-A 398 408 10 20 231 HULA-A 107 10 15 194 HULA-A 107 10 15 194 HULA-A 107 10 15 194 HULA-A 300 785 515 101 HULA-A 107 10 15 194 HULA-A 107
168 HULA-A 175 185 10 15 21 HULA-A 212 229 17 20 181 HULA-A 260 265 5 15 178 HULA-A 260 265 5 15 192 HULA-A 315 320 5 15 194 HULA-A 340 350 10 15 194 HULA-A 365 375 10 15 38 HULA-A 398 408 10 20 43 HULA-A 451 467 16 20 210 HULA-A 455 570 75 15 217 HULA-A 585 590 5 15 65 HULA-A 676 691 15 20 231 HULA-A 710 720 10 15 237 HULA-A 780 785 5 15 86 HULA-A 869 884 15 20 252 HULA-A 915 920 5 15
252 HULA-A 915 920 5 15 249 HULA-A 935 940 5 15 256 HULA-A 955 985 30 15 101 HULA-A 1022 1038 16 20 271 HULA-A 1071 1086 15 15 112 HULA-A 1124 1137 13 20 280 HULA-A 1124 1137 13 20 280 HULA-A 1200 1205 5 15 287 HULA-A 1260 1275 15 15 288 HULA-A 1290 1295 5 15 302 HULA-A 1345 1359 14 15 147 HULA-B 0 10 15 153 HULA-B 30 50 20 15 165 HULA-B 145 145 20 15
212 HULA-B 565 605 40 15 220 HULA-B 565 605 40 15 230 HULA-B 705 750 45 15 76 HULA-B 778 792 14 20 96 HULA-B 991 1007 16 20 102 HULA-B 1034 1048 14 20 275 HULA-B 1105 1120 15 15 281 HULA-B 1175 1184 9 15 124 HULA-B 1245 1258 13 20 298 HULA-B 1285 1325 40 15
4 COUX-A 20 34 14 20 167 CDOK-A 175 220 45 15 29 CDOK-A 315 330 15 20 197 CDOK-A 375 390 15 15 204 CDOK-A 375 390 15 15 204 CDOK-A 420 435 15 15 204 CDOK-A 420 435 15 15 204 CDOK-A 420 435 15 15 208 COOK-A 520 5.05 10 15 208 COOK-A 540 565 25 15 222 COOK-A 580 590 10 15 239 COOK-A 690 705 15 15 239 COOK-A 858 874 16 20 92 COOK-A 917 933 16 20

261 COOKK- 268 COOKK- 268 COOKK- 268 COOKK- 268 COOKK- 268 COOKK- 269 COOKK- 260 COOKK-
242 1101 247 1101 255 1101 255 1101 267 1101 267 1101 276 1101 276 1101 284 1101 301 1101 309 1101 309 1101 311 1101 309 1101 311 1105 9 BOSS-7 13 BOSS
265 BOSS 278 BOSS
119 BOSS 283 BOSS 293 BOSS 1294 BOSS 144 BOSS 162 BOSS 162 BOSS 177 BOSS 205 BOSS 205 BOSS 205 BOSS 205 BOSS 205 BOSS 2246 BOSS 2246 BOSS 2246 BOSS 2246 BOSS 2246 BOSS 2246 BOSS 2246 BOSS 200

OOK-A	990 1010 20 15 1051 1061 10 15 1089 1103 14 20
OOK-A OOK-A	1152 1166 14 20 1190 1200 10 15 1267 1307 40 15 1395 1408 13 20
00K-A 00K-B 00K-B	1395 1408 13 20 181 197 16 20 265 740 475 15
00K-B 00K-B	333 847 14 20 980 1170 190 15 1260 1285 25 15
OOK-B	1260 1285 25 15 1315 1329 14 20 1391 1404 13 20 30 35 5 15
IDI-A IDI-A	52 62 10 15 80 90 10 15 175 180 5 15
IDI-A	200 220 20 15 265 285 20 15 325 285 10 15
	345 365 20 15 420 425 5 15
IDI-A	520 535 15 20 560 570 10 15
	630 640 10 15 655 665 10 15
	800 815 15 15 830 850 20 15
	935 955 20 15 977 991 14 20
IDI-A IDI-A	1006 1011 5 15 1040 1050 10 15 1070 1085 15 15
INDI-A INDI-A INDI-A	1110 1135 25 15 1208 1218 10 15 1272 1288 16 20
INDI-A INDI-A INDI-A	1320 1330 10 151370 1376 5 151395 1405 10 151420 1440 20 15
INDI-A BOSS-A SS-A 1	1420 1440 20 15 15 60 45 15 00 115 15 20
DSS-A BOSS-A DSS-A	990 1010 20 15 1061 1061 10 15 1089 1103 14 20 1152 1166 14 20 1190 1200 10 15 1267 1307 40 15 1395 1404 120 181 197 16 20 265 740 475 15 33 847 14 20 980 1170 190 15 1260 1285 25 15 1315 1329 14 20 30 35 5 15 52 62 10 15 80 90 10 15 175 180 5 15 266 285 20 15 326 330 10 15 345 365 20 15 326 470 20 15 345 365 20 15 320 535 15 20 560 570 10 15 630 640 10 15 630 640 10 15 630 640 10 15 630 850 20 15 800 815 15 15 800 815 15 15 800 815 15 15 1040 1050 10 15 1770 1085 15 15 1040 1050 10 15 1701 1376 5 15 1040 1050 10 15 1770 1376 5 15 1040 1050 10 15 1370 1376 5 15 1395 1406 10 15 1420 144 20 15 379 396 17 20 745 755 10 15 800 805 20 15 15 20 15 20
BOSS-A DSS-A DSS-A	308 328 20 15 379 396 17 20 483 497 14 20
)\$\$-A)\$\$-A)\$\$-A	535 549 14 20 582 594 12 20 634 649 15 20
ĴŠŠ-Ă ĴSS-A ROSS-A	679 693 14 20 713 730 17 20 745 755 10 15
BOSS-Â BOSS-A	770 790 20 15 850 885 35 15 959 971 12 20
DSS-A DSS-A BOSS-A BOSS-A	850 885 35 15 959 971 12 20 995 995 0 20 1020 1040 20 15 1125 1130 5 15
BUSS-A	$\begin{array}{cccccccccccccccccccccccccccccccccccc$
BOSS-A BOSS-A BOSS-A BOSS-A BOSS-A BOSS-A BOSS-B BOSS-B BOSS-B	1293 1307 14 20 1393 1409 16 20 1424 1444 20 15
BOSS-B BOSS-B BOSS-B BOSS-B	$\begin{array}{c} 1020 & 1040 & 20 & 15 \\ 1125 & 1130 & 5 & 15 \\ 1196 & 1210 & 14 & 20 \\ 1230 & 1240 & 10 & 15 \\ 1260 & 1265 & 5 & 15 \\ 1293 & 1307 & 14 & 20 \\ 1393 & 1409 & 16 & 20 \\ 1424 & 1444 & 20 & 15 \\ 115 & 190 & 75 & 15 \\ 205 & 215 & 10 & 15 \\ 230 & 280 & 50 & 15 \\ 345 & 390 & 45 & 15 \\ 490 & 505 & 15 & 15 \\ 555 & 570 & 15 \\ 555 & 570 & 15 \\ 555 & 570 &$
BOSS-B BOSS-B	115 190 75 15 205 215 10 15 230 280 50 15 345 390 45 15 490 505 15 15 555 570 15 15
BOSS-B DSS-B BOSS-B	555 570 15 15 613 627 14 20 642 662 20 15 685 720 35 15 860 865 5 15
BOSS-B BOSS-B DSS-B	555 570 15 16 613 627 14 20 642 662 20 15 685 720 35 15 860 865 5 15 885 900 15 20
BOSS-A BOSS-A BOSS-A BOSS-A BOSS-A BOSS-B BO	
BOSS-B	1280 1290 10 15

141 308 161	BOSS-B BOSS-B LICE-A	1369 1394 85 10	1379 1414 5 20	10 20 20 15 15
10 23 27 191	ION-A ION-A LION-A	191 20 239 25 308 32 337 3	$ \begin{array}{c} 14 \\ 3 \\ 14 \\ 2 \\ 14 \\ 47 \\ 10 \\ 1 $	20 20 15
203 48 L 228		420 4 511 52 690 7	80 60 4 13 10 20	15 20 15
91 Î 98 L 262 107	ÍON-A ION-A LION-A	904 91 997 10 1026 1094	9 15 11 14 1046 1108	20 20 15 14 20
118 130 143 150	LION-A LION-A LION-A LION-B	1194 1298 1393 20 25	1210 1311 1402 5 15	14 20 16 20 13 20 9 20
157 173 31 I 200	LION-B LION-B ION-B LION-B	60 18 195 2 336 35 390 4	0 120 10 15 0 14 20 30	15 15 20 15
207 241 90 I 99 I	LION-B LION-B ION-B ION-B	480 5 810 8 900 91 1002 1	25 45 45 35 3 13 016 1	15 15 20 4_20_
272 300 305 15 0	LION-B LION-B LION-B UAM-A	1080 1300 1365 171 18	$ \begin{array}{r} 1100 \\ 1310 \\ 1405 \\ 3 12 \\ \hline \end{array} $	20 15 10 15 40 15 20
20 28 39 53 63	UAN-A UAN-A UAN-A	212 22 308 31 403 41 578 59	9 11 7 14 4 16	20 20 20 20
77 94 105 274	UAM-A UAM-A GUAM-A GUAM-A	797 81 976 99 1064 1105	0 13 1 15 1078 1200	20 20 14 20 95 15
122 291 299 135	GUAM-A GUAM-A GUAM-A GUAM-A	1226 1260 1295 1327	1241 1280 1305 1342	95 15 15 20 20 15 10 15 15 20
304 235 103 296	GUAM-A GUAM-B GUAM-B GUAM-B	1357 725 9 1036 1275	1377 15 19 1046 1290	10 15 15 20 20 15 00 15 10 20 15 15
3 P] 6 P] 172 54 F	IKE-A 1 IKE-A 8 PIKE-A PIKE-A	9 30 1 2 97 1 180 3 581 59	1 20 5 20 60 18	14 20 10 15
61 F 66 F 73 F 81 F	PIKE-A PIKE-A PIKE-A PIKE-A	635 64 678 69 732 74 815 83	8 13 2 14 6 14 2 17	20 20 20 20
263 110 115 120	PIKE-A PIKE-A PIKE-A PIKE-A	990 1 1112 1154 1212	035 4 1126 1169 1228	5 15 14 20 15 20 16 20 45 20
134 11 F 22 F 187	FIAD-A LEEF-A REEF-A REEF-A	123 13 123 13 226 23 300 3 357 37	1305 19 16 17 11 120 20	45 20 20 20 15 20
40 F 213 59 F 71 F	BBCLION - A A A A A A A A A A A A A A A A A A	422 43 515 5 621 63 719 73	2^{15} 3^{11} 3^{5} 1^{10} 1^{12	20 15 20 20
74 F 87 F 116 126 137		875 89 1171		20 15 20 15 20
137 146	REEF-A REEF-A	1335 1431	1347 1445	12 20 14 20

Day 4

284 POGO-C 1110 1155 45 15 287 POGO-C 1170 1185 15 15 293 POGO-C 1200 1215 15 15
287 POGO-C 1170 1185 15 15 293 POGO-C 1200 1215 15 15 120 POGO-C 1247 1263 16 20 127 POGO-C 1305 1320 15 20
120 POGO-C 1247 1263 16 20 127 POGO-C 1305 1320 15 20 134 POGO-C 1369 1382 13 20 139 POGO-C 1407 1420 13 20 149 HULA-A 25 40 15 15
120 POGO-C 1247 1263 16 20 127 POGO-C 1306 1320 15 20 134 POGO-C 1369 1362 13 20 139 POGO-C 1369 1362 13 20 139 POGO-C 1407 1420 13 20 149 HULA-A 25 40 15 15 155 HULA-A 60 70 10 15 150 HULA-A 85 105 20 15 10 HULA-A 85 105 20 15 10 HULA-A 199 216 17 20 17 HULA-A 199 216 17 20
150 HULA-A 85 105 20 15 10 HULA-A 142 157 15 20 17 HULA-A 199 216 17 20 167 HULA-A 231 236 5 15 174 HULA-A 251 271 20 15
167 HULA-A 199 216 1/ 20 167 HULA-A 231 236 5 15 174 HULA-A 251 271 20 15 181 HULA-A 290 295 5 15 188 HULA-A 345 390 45 15 38 HULA-A 419 430 11 20 202 HULA-A 419 430 11 20 202 HULA-A 420 500 10 15 48 WULA-A 420 500 10 15
120 POGO-C 1247 1263 16 20 127 POGO-C 1305 1320 15 20 134 POGO-C 1369 1382 13 20 139 POGO-C 1407 1420 13 20 149 HULA-A 25 40 15 15 155 HULA-A 85 105 20 15 10 HULA-A 85 105 20 15 10 HULA-A 142 157 15 20 17 HULA-A 231 236 5 15 174 HULA-A 290 295 5 15 188 HULA-A 2490 295 5 15 188 HULA-A 419 430 11 20 202 HULA-A 490 500 10 15 48 HULA-A 520 500 10 15 48 HULA-A 500 500 10 15 58 HULA-A 500 500 500 500 500 500 500 500 500 50
211 HULA-A 615 635 20 15 225 HULA-A 660 670 10 15 232 HULA-A 705 720 15 15
48 HULA-A 522 533 11 20 218 HULA-A 580 600 20 15 211 HULA-A 615 635 20 15 225 HULA-A 660 670 10 15 232 HULA-A 705 720 15 15 233 HULA-A 705 720 15 15 78 HULA-A 804 818 14 20 247 HULA-A 850 925 75 15 258 HULA-A 940 945 5 15 95 HULA-A 1032 1047 15 20 107 HULA-A 1079 1091 12 20 286 HULA-A 1125 1130 5 15 291 HULA-A 1185 1190 5 15 296 HULA-A 1230 1275 45 15 296 HULA-A 1230 1275 45 15
258 HULA-A 940 945 5 15 95 HULA-A 976 993 17 20 99 HULA-A 1032 1047 15 20
99 HULA-A 1032 1047 15 20 107 HULA-A 1079 1091 12 20 286 HULA-A 1125 1130 5 15 291 HULA-A 1185 1190 5 15 296 HULA-A 1230 1275 45 15 302 HULA-A 1290 1330 40 15 136 HULA-A 1377 1390 13 20 309 HULA-A 1405 1415 10 15 318 HULA-A 1405 1415 15 15 158 HULA-B 70 80 10 15 164 HULA-B 74 80 10 15 172 HULA-B 225 230 5 15
200 HULA-A 1125 1130 5 15 291 HULA-A 1185 1190 5 15 302 HULA-A 1230 1275 45 15 302 HULA-A 1290 1330 40 15 136 HULA-A 1377 1390 13 20 309 HULA-A 1405 1415 10 15 318 HULA-A 1430 1445 15 15 158 HULA-B 70 80 10 15 154 HULA-B 70 80 10 15
136 HULA-A 1377 1390 13 20 309 HULA-A 1405 1415 10 15 318 HULA-A 1430 1445 15 15 158 HULA-B 70.80 1445 15 15
164 HULA-B 140 160 20 15 172 HULA-B 225 230 5 15 170 HULA-B 245 255 10 15
170 HULA-B 245 255 10 15 180 HULA-B 285 305 20 15 187 HULA-B 340 350 10 15 191 HULA-B 385 395 10 15 191 HULA-B 385 395 10 15
187 HULA-B 340 350 10 15 191 HULA-B 385 395 10 15 37 HULA-B 419 436 17 20 47 HULA-B 515 526 11 20 213 HULA-B 515 526 11 20 213 HULA-B 565 605 40 15 228 HULA-B 720 735 15 15 235 HULA-B 720 735 15 15 241 HULA-B 820 825 5 15 83 HULA-B 960 9915 13 20 261 HULA-B 950 960 10 15 98 HULA-B 994 1009 15 20 266 HULA-B 1095 1110 15 20 289 HULA-B 1095 1110 15 20 289 HULA-B 1155 1180 25 15 123 HULA-B 1275 1290 15 20 300 HULA-B 1345 1360 15 15 307 HULA-B 1345 1360 15 15 4 COOK-A 52 65 13 20 165 COX-A 160 165 5 15
213 HULA-B 565 605 40 15 228 HULA-B 680 695 15 15 235 HULA-B 720 735 15 15 241 HULA-B 720 735 15 15 83 HULA-B 820 825 5 15 83 HULA-B 902 915 13 20 90 HULA-B 902 915 13 20
241 HULA-B 820 825 5 15 83 HULA-B 866 880 14 20 90 HULA-B 902 915 13 20 261 HULA-B 950 960 10 15 98 HULA-B 994 1009 15 20 266 HULA-B 1024 1044 20 15 109 HULA-B 1095 1110 15 20 289 HULA-B 1155 1180 25 15 294 HULA-B 1200 1245 45 15
98 HULA-B 994 1009 15 20 266 HULA-B 1024 1044 20 15 109 HULA-B 1095 1110 15 20
261 HULA-B 950 960 10 15 98 HULA-B 994 1009 15 20 266 HULA-B 1024 1044 20 15 109 HULA-B 1095 1110 15 20 289 HULA-B 1155 1180 25 15 294 HULA-B 1200 1245 45 15 123 HULA-B 1275 1290 15 20
300 HULA-B 1305 1320 15 15 307 HULA-B 1345 1360 15 15 4 COOK-A 52 65 13 20 165 COOK-A 160 165 5 15
284 POGO-C 1110 1155 45 15 287 POGO-C 1200 1215 15 15 120 POGO-C 1207 1263 16 20 127 POGO-C 1369 1362 13 20 139 POGO-C 1369 1362 13 20 149 HULA-A 25 40 15 15 150 HULA-A 60 70 10 15 150 HULA-A 142 157 15 20 17 HULA-A 291 236 5 15 174 HULA-A 290 295 5 15 188 HULA-A 345 390 45 15 38 HULA-A 490 500 10 15 181 HULA-A 520 533 11 20 202 HULA-A 490 500 10 15 232 HULA-A 660 670 10 15 233 HULA-A 550 600 20 16 211 HULA-A 550 600 20 16 211 HULA-A 550 600 20 15 233 HULA-A 660 670 10 15 233 HULA-A 705 720 15 15 233 HULA-A 705 720 15 15 233 HULA-A 1032 1047 15 20 107 HULA-A 1125 1130 5 15 233 HULA-A 1125 1130 5 15 296 HULA-A 1290 1330 40 15 136 HULA-A 1079 1091 12 20 286 HULA-A 1230 1275 45 15 396 HULA-A 1079 1091 12 20 299 HULA-A 1079 1091 12 20 286 HULA-A 1230 130 40 15 136 HULA-A 1377 1390 13 20 107 HULA-A 1230 1275 45 15 302 HULA-A 1405 110 5 15 291 HULA-A 1405 110 5 15 292 HULA-A 1405 1130 5 15 291 HULA-A 1290 1330 40 15 136 HULA-A 1290 1330 40 15 136 HULA-A 1405 10 15 318 HULA-A 1405 10 15 318 HULA-B 70 80 10 15 172 HULA-B 245 255 10 15 172 HULA-B 140 160 20 15 174 HULA-B 140 160 20 15 174 HULA-B 140 160 20 15 175 HULA-B 140 160 20 15 174 HULA-B 140 160 20 15 175 HULA-B 140 160 20 15 174 HULA-B 140 160 20 15 175 HULA-B 140 160 20 15 174 HULA-B 140 160 20 15 175 HULA-B 140 160 20 15 174 HULA-B 140 160 20 15 175 HULA-B 140 160 20 15 174 HULA-B 140 160 20 15 175 HULA-B 140 160 20 15 174 HULA-B 140 160 20 15 175 HULA-B 140 160 20 15 174 HULA-B 140 160 20 15 175 HULA-B 140 160 20 15 174 HULA-B 140 160 20 15 175 HULA-B 140 160 20 15 174 HULA-B 155 1526 11 20 215 HULA-B 155 1526 15 235 HULA-B 1005 110 15 20 266 HULA-B 1005 110 15 20 260 HULA-B 10
189 COOK-A 378 393 15 15 209 COOK-A 540 550 10 15 219 COOK-A 585 600 15 15 229 COOK-A 680 715 35 15 229 COOK-A 680 715 35 15
229 COOK-A 680 715 35 15 244 COOK-A 840 850 10 15
268 CODK-A 990 1035 45 15 103 CDDK-A 1058 1073 15 20 283 CDDK-A 1105 1120 15 15 297 CDK-A 1245 1260 15 15 125 CDCK-A 1245 1360 16 20
268 COOK-A 990 1035 45 15 103 COOK-A 1058 1073 15 20 283 COOK-A 1058 1073 15 20 283 COOK-A 1105 1120 15 15 125 COOK-A 1245 1260 15 15 125 COOK-A 1245 1260 15 15 125 COOK-A 1284 1300 16 20 141 COOK-A 1426 1440 14 20 3 COOK-B 48 61 13 20 173 COOK-B 225 275 50 15 193 COOK-B 390 415 25 15 207 COOK-B 390 415 25 15 207 COOK-B 570 580 15 15 217 COOK-B 570 580 15 15 226 COOK-B 660 675 15 15
141 COOK-A 1420 144 20 3 COOK-B 48 61 13 20 173 COOK-B 225 275 50 15 193 COOK-B 390 415 25 15 207 COOK-B 530 545 15 15 217 COOK-B 570 580 10 15 226 COOK-B 570 580 10 15 227 COOK-B 570 580 570 580 570 580 50 228 COOK-B 570 580 580 570 580 50 228 COOK-B 570 580 580 580 580 580 580 580 580 580 58
207 COOK-B 530 545 15 15 217 COOK-B 570 580 10 15 226 COOK-B 660 675 15 15

248 COOK-B 850 865 15 15 253 COOK-B 900 945 45 15 262 COOK-B 960 970 10 15 267 COOK-B 990 1655 665 15 144 UBU-A 30 35 5 15
248 COOK-B 850 865 15 15 253 COOK-B 900 945 45 15 262 COOK-B 960 970 10 15 267 COOK-B 990 1655 665 15 144 IWDI-A 53 63 10 15 161 IWDI-A 90 100 10 15 159 IWDI-A 115 135 20 15 168 IWDI-A 180 185 5 15 171 IWDI-A 220 265 45 15 185 IWDI-A 320 330 10 15 185 IWDI-A 345 365 20 15
186 INDI-A 345 365 20 15 192 INDI-A 390 416 25 15 198 INDI-A 450 470 20 15 203 INDI-A 500 505 5 15 206 INDI-A 500 540 20 15 212 INDI-A 560 570 10 15 210 INDI-A 585 595 10 15 224 INDI-A 585 595 10 15
223 ÎNDÎ-A 680 700 20 15 68 INDI-A 720 735 15 20 239 INDI-A 800 810 10 15 242 ÎNDI-A 830 850 20 15 251 ÎNDI-A 870 890 20 15 250 INDI-A 905 925 20 15 94 INDI-A 948 964 16 20 259 INDI-A 979 989 10 15 273 INDI-A 979 989 10 15
272 INDI-A 1040 1050 10 15 276 INDI-A 1065 1080 15 15 282 INDI-A 1105 1200 95 15 292 INDI-A 1215 1235 20 15 293 INDI-A 1215 1235 20 15
250 1101-A 1295 1305 10 15 131 IWDI-A 1344 1357 13 20 137 IWDI-A 1344 1357 13 20 137 IWDI-A 1344 1357 13 20 137 IWDI-A 1347 1399 12 20 315 IWDI-A 1387 1399 12 20 315 IWDI-A 1414 1424 10 15 311 IWDI-A 1439 1444 5 15 5 BOSS-A 54 70 16 20 157 BOSS-A 85 90 5 15 8 BOSS-A 125 140 15 20
21 BOSS-A 252 262 10 20 177 BOSS-A 277 297 20 15 30 BOSS-A 349 365 16 20 36 BOSS-A 419 431 12 20 43 BOSS-A 469 481 12 20 204 BOSS-A 510 540 30 15 55 BOSS-A 584 597 13 20 214 BOSS-A 612 632 20 15
214 BOSS-A 667 682 15 20 236 BOSS-A 750 770 20 15 245 BOSS-A 840 850 10 15 252 BOSS-A 885 890 5 15 97 BOSS-A 979 994 15 20 271 BOSS-A 1020 1040 20 15 275 BOSS-A 1055 1065 10 15 279 BOSS-A 1080 1090 10 15 281 BOSS-A 1105 1120 15 15 117 BOSS-A 1226 1241 15 20
279 BOSS-A 1080 1090 10 15 281 BOSS-A 1105 1120 15 15 117 BOSS-A 1226 1241 15 20 303 BOSS-A 1285 1295 10 15 128 BOSS-A 1285 1295 10 15 128 BOSS-A 1326 1342 16 20 314 BOSS-A 1380 1385 5 15 142 BOSS-A 1434 1448 14 20
142 B053-B 1434 40 16 20 156 B055-B 440 16 20 156 B055-B 425 440 15 15 200 B055-B 425 440 15 15 200 B055-B 480 485 5 15 54 B055-B 576 592 16 20 220 B055-B 639 656 17 20 67 B055-B 639 656 17 20 67 B055-B 685 701 16 20 237 B055-B 760 775 15 15 243 B055-B 1366 1380 14 20 312 B055-B 1366 1380 14 20 312 B055-B 1395 1415 20 15 145 L10N-A 10 15 5 15 162 L10N-A 269 284 15 20 34 L10N-A 269 284 15 20 34 L10N-A 450 455 5 15 44 L10N-A 482 496 14 20
196 BUSS-B 425 440 15 15 200 BUSS-B 480 485 5 15 54 BUSS-B 576 592 16 20 220 BUSS-B 639 656 17 20 67 BUSS-B 639 656 17 20 67 BUSS-B 685 701 16 20 237 BUSS-B 760 775 15 15 243 BUSS-B 760 775 15 15 133 BUSS-B 1366 1380 14 20 312 BUSS-B 1366 1380 14 20 312 BUSS-B 1395 1415 20 15 145 LION-A 10 15 5 15 162 LION-A 90 115 25 15 18 LION-A 223 237 14 20 25 LION-A 269 284 15 20 34 LION-A 380 396 16 20
25 LION-A 269 284 15 20 34 LION-A 380 396 16 20 197 LION-A 450 455 5 15 44 LION-A 482 496 14 20 208 LION-A 540 565 25 15 222 LION-A 605 625 20 15 66 LION-A 675 689 14 20

Day 5

00 D000 4 070 004 40 00
86 POGD-A 879 891 12 20 87 POGD-A 911 922 11 20 92 POGD-A 947 960 13 20 95 POGD-A 947 960 13 20
92 POGO-A 947 960 13 20 253 POGO-A 975 990 15 15 261 POGO-A 1005 1020 15 15 103 POGO-A 1048 1061 13 20
270 POGO-A 1100 1015 15 15 273 POGO-A 1130 1140 10 15 276 POGO-A 1165 1175 10 15
276 POGO-A 1165 1175 10 15 120 POGO-A 1211 1227 16 20 285 POGO-A 1255 1280 25 15
103 POGO-A 1048 1061 13 20 270 POGO-A 1100 1115 15 15 273 POGO-A 1130 1140 10 15 276 POGO-A 1165 1175 10 15 120 POGO-A 1255 1280 25 15 129 POGO-A 1255 1280 25 15 129 POGO-A 1310 1326 16 20 295 POGO-A 1350 1365 15 15 139 POGO-A 1410 1426 16 20
143 POGO-B 0 10 10 15 A POGO-B 35 51 16 20
276 POGO-A 1165 1175 10 15 120 POGO-A 1211 1227 16 20 285 POGO-A 1211 12280 25 15 129 POGO-A 1310 1326 16 20 295 POGO-A 1310 1426 16 20 143 POGO-B 0 10 15 4 POGO-B 96 101 15 156 POGO-B 96 101 5 156 POGO-B 129 143 14 20 20 POGO-B 167 181 14 20 20 POGO-B 263 278 15 20 176 POGO-B 283 303 10 15 30 POGO-B 328
11 P0G0-B 129 143 14 20 16 P0G0-B 167 181 14 20 20 P0G0-B 204 217 13 20 23 P0G0-B 263 278 15 20
176 POGD-B 293 303 10 15 30 POGD-B 328 344 16 20 36 POGD-B 391 404 13 20 40 POGD-B 429 445 16 20
36 PUGU-B 391 404 13 20 40 PUGU-B 429 445 16 20 46 PUGU-B 489 503 14 20
46 p0G0-B 489 503 14 20 198 P0G0-B 518 523 5 15 52 p0G0-B 545 562 17 20 56 P0G0-B 587 602 16 20 62 P0G0-B 683 649 16 20 68 P0G0-B 683 649 16 20 68 P0G0-B 883 649 16 20 68 P0G0-B 883 649 16 20 68 P0G0-B 883 854 16 20 82 P0G0-B 883 854 16 20 240 P0G0-B 869 899 30 15 251 P0G0-B 914 929 15 15
62 POGO-B 633 649 16 20 68 POGO-B 684 699 15 20
77 PUGU-B 781 795 14 20 82 PUGU-B 838 854 16 20 240 PUGU-B 869 899 30 15
251 POGO-B 914 929 15 15 95 POGO-B 955 966 11 20 260 POGO-B 990 1000 10 15
100 POGO-B 1020 1036 16 20 107 POGO-B 1071 1086 15 20 109 POGO-B 1120 1137 17 20
109 P0G0-B 1120 1137 17 20 275 P0G0-B 1152 1162 10 15 115 P0G0-B 1189 1202 13 20
95 POGO-B 955 966 11 20 260 POGO-B 995 966 11 20 100 POGO-B 990 1000 10 15 100 POGO-B 1020 1036 16 20 107 POGO-B 1071 1086 15 20 109 POGO-B 1152 1162 10 15 115 POGO-B 1189 1202 13 20 123 POGO-B 1246 1260 14 20 127 POGO-B 1246 1260 14 20 127 POGO-B 1290 1305 15 20 292 POGO-B 1335 1350 15 15 137 POGO-B 1392 1408 16 20 304 POGO-B 1423 1433 10 15
137 POGO-B 1392 1408 16 20 304 POGO-B 1423 1433 10 15 6 POGO-C 67 84 17 20
10 POGO-B 328 344 16 20 36 POGO-B 391 404 13 20 40 POGO-B 391 404 13 20 46 POGO-B 489 503 14 20 198 POGO-B 649 562 17 20 56 POGO-B 687 602 15 20 62 POGO-B 687 602 15 20 62 POGO-B 688 699 15 20 70 POGO-B 838 854 16 20 70 POGO-B 914 929 15 15 75 POGO-B 914 929 15 15 76 POGO-B 990 1000 10 15 70 POGO-B 990 1000 10 15 70 POGO-B 152 1162 10 15 70 POGO-B 120 1036 15 20 107
189 PUGU-C 220 275 55 15 182 PUGU-C 330 345 15 15 190 PUGU-C 375 395 20 15
160 $\overline{P}0G0-C$ 105 120 15 15 17 $\overline{P}0G0-C$ 168 184 16 20 169 $\overline{P}0G0-C$ 220 275 55 15 182 $\overline{P}0G0-C$ 330 345 15 15 190 $\overline{P}0G0-C$ 375 395 20 15 195 $\overline{P}0G0-C$ 420 435 15 15 199 $\overline{P}0G0-C$ 420 435 15 15 199 $\overline{P}0G0-C$ 450 460 10 15 199 $\overline{P}0G0-C$ 480 495 15 15 51 $\overline{P}0G0-C$ 580 595 15 15 63 $\overline{P}0G0-C$ 642 658 16 20 228 $\overline{P}0G0-C$ 642 658 16 20 228 $\overline{P}0G0-C$ 714 728 14 20 73 $\overline{P}0G0-C$ 748 764 16 20
51 P0G0-C 531 547 16 20 220 P0G0-C 580 595 15 15 43 P0G0-C 642 658 16 20
228 P0G0-C 673 683 10 15 70 P0G0-C 714 728 14 20
73 PUGD-C 748 764 16 20 243 PUGD-C 845 860 15 15 88 PUGD-C 914 929 15 20
88 P060-C 914 929 15 20 93 P060-C 949 962 13 20 254 P060-C 977 1007 30 15 267 P060-C 1050 1060 10 15
113 POGO-C 1152 1166 14 20 125 POGO-C 1256 1268 12 20 289 POGO-C 1283 1323 40 15 133 POGO-C 1346 1361 15 20 301 POGO-C 1395 1430 35 15 149 HULA-A 30 50 20 15
289 POGO-C 1283 1323 40 15 133 POGO-C 1346 1361 15 20 301 POGO-C 1395 1430 35 15
167 VIII Ama 136 166 7() 16
170 HIII.A-A 240 260 20 15
171 HULA-A 310 325 15 15 185 HULA-A 340 350 10 15
35 HULA-A 389 404 15 20 39 HULA-A 426 437 11 20 202 HULA-A 485 495 10 15
202 HULA-A 485 495 10 15 204 HULA-A 510 530 20 15 215 HULA-A 570 580 10 15

216 HULA-A 595 605 10 15 224 HULA-A 650 635 5 15 225 HULA-A 765 800 35 15 81 HULA-A 765 800 35 15 81 HULA-A 895 940 45 15 97 HULA-A 962 978 16 20 249 HULA-A 962 978 16 20 256 HULA-A 1000 1014 14 20 256 HULA-A 1000 1014 14 20 256 HULA-A 1110 115 5 15 269 HULA-A 1110 115 5 15 269 HULA-A 11205 1218 13 20 281 HULA-A 11205 1218 13 20 281 HULA-A 1305 1320 15 20 135 HULA-A 1389 1399 10 15 151 HULA-B 1358 1374 16 20 297 HULA-B 1389 1399 10 15 151 HULA-B 1355 15 165 HULA-B 1355 15 165 HULA-B 1355 15 165 HULA-B 135 1320 15 20 135 HULA-B 135 1320 15 20 135 HULA-B 135 1320 15 15 165 HULA-B 135 1320 15 15 161 HULA-B 135 1320 15 15 161 HULA-B 135 1320 15 15 165 HULA-B 180 195 15 15 26 HULA-B 180 195 15 15 26 HULA-B 755 615 40 15 177 HULA-B 489 503 14 20 207 HULA-B 575 615 40 15 214 HULA-B 775 810 35 15 233 HULA-B 775 810 35 15 244 HULA-B 190 1200 10 15 227 HULA-B 190 1200 10 15 233 HULA-B 190 1200 10 15 233 HULA-B 190 1200 10 15 244 HULA-B 489 503 14 20 207 HULA-B 190 1200 15 214 HULA-B 190 1200 10 15 227 HULA-B 489 503 12 238 HULA-B 775 810 35 15 244 HULA-B 1190 1200 10 15 238 HULA-B 1190 1200 10 15 238 HULA-B 1190 1200 10 15 244 HULA-B 1290 1295 5 15 244 HULA-B 1290 1295 5 15 244 HULA-B 1290 1295 5 15 245 HULA-B 1290 1295 5 15 246 CO0K-A 70 75 5 15 177 CO0K-A 280 2975 13 20 262 CO0K-A 70 55 15 174 CO0K-A 280 2975 13 20 262 CO0K-A 1010 1025 15 15 277 CO0K-A 1356 1367 11 20 173 CO0K-A 1356 1367 11 20 174 CO0K-B 885 930 45 15 154 CO0K-A 1356 1367 11 20 155 1MD1-A 135 1286 50 15 157 1MD1-A 53 63 10 15 158 1MD1-A 1356 1260 45 15 247 CO0K-A 1356 1367 11 20 152 1MD1-A 135 1260 45 15 247 CO0K-A 1356 1367 11 20 152 1MD1-A 135 120 5 15 158 1MD1-A 135 120 5 15 159 1MD1-A 135 120 5 15 159 1MD1-A 135 120 5 15	
200 COOK-A 700 756 15 15 218 COOK-A 570 536 15 15 218 COOK-A 645 650 5 15 78 COOK-A 794 808 14 20 84 COOK-A 962 975 13 20 262 COOK-A 1010 1025 15 15 124 COOK-A 1010 1025 15 15 124 COOK-A 1254 1270 16 20 287 COOK-A 1285 1290 5 15 134 COOK-A 1356 1367 11 20 173 COOK-B 260 730 470 15 241 COOK-B 825 870 45 15 263 COOK-B 1235 1285 50 15 136 COOK-B 1235 1285 50 15 136 COOK-B 1235 1285 50 15 157 INDI-A 53 63 10 15 158 INDI-A 105 110 5 15 163 INDI-A 135 140 5 15 158 INDI-A 136 182 16 20 167 INDI-A 235 255 20 15 179 INDI-A 236 255 20 15 181 INDI-A 330 335 5 15 183 INDI-A 375 395 20 15 233 INDI-A 417 430 13 20 197 INDI-A 57 650 10 15 233 INDI-A 40 470 20 15 233 INDI-A 40 470 20 15 233 INDI-A 40 655 10 5 15 241 COOK 55 10 5 15 250 INDI-A 670 580 10 15 231 INDI-A 80 90 10 15 231 INDI-A 80 90 10 15 233 INDI-A 417 430 13 20 197 INDI-A 80 696 5 15 230 INDI-A 710 730 20 15 239 INDI-A 80 810 10 15 239 INDI-A 80 810 10 15 239 INDI-A 80 810 10 15 245 INDI-A 80 810 10 15 250 INDI-A 80 810 10 15 250 INDI-A 80 810 10 15 250 INDI-A 80 810 10 15 251 INDI-A 80 810 10 15 252 INDI-A 80 810 10 15 252 INDI-A 80 810 10 15 255 INDI-A 80 80 90 10 15 255 INDI-A 80 80 90 10 15 255 INDI-A 80 80 90 10 15 255 INDI-A 80 810 10 15 255 INDI-A 905 910 5 15 255 INDI-	

193 GUAM-B 405 450 45 15 110 GUAM-B 1128 1141 13 20 2 PIKE-A 27 38 11 20 10 PIKE-A 125 141 16 20 27 PIKE-A 306 321 15 20
10 PIKE-A 125 141 16 20 27 PIKE-A 306 321 15 20 53 PIKE-A 548 560 12 20 57 PIKE-A 599 611 12 20 64 PIKE-A 644 658 14 20 69 PIKE-A 645 710 15 20
75 PIKE-A 758 774 16 20 112 PIKE-A 1152 1168 16 20 303 PIKE-A 1410 1430 20 15 147 REF-A 10 30 20 15
5 REEF-A 67 81 14 20 43 REEF-A 460 476 16 20 67 REEF-A 684 698 14 20 80 REEF-A 818 834 16 20
98 REEF-A 975 993 18 20 268 REEF-A 1080 1100 20 15 119 REEF-A 1211 1227 16 20
132 REEF-A 1320 1332 12 20 138 REEF-A 1394 1409 15 20 305 REEF-A 1425 1435 10 15

254 POGO-B 1045 1055 10 15 103 POGO-B 1086 1099 13 20 107 POGO-B 1150 1165 15 15 120 POGO-B 1130 1195 15 15 120 POGO-B 1234 1246 12 20 134 POGO-B 1353 1367 14 20 134 POGO-C 130 150 20 15 161 POGO-C 130 150 20 15 161 POGO-C 130 150 20 15 188 POGO-C 130 20 15 188 POGO-C 130 445 15 15 188 POGO-C 399 415 16 20 188 POGO-C 430 445 15 15 197 POGO-C 588 602 14 20 211 POGO-C 617 632 15 15 188 POGO-C 782 797 15 20 231 POGO-C 845 865 20 15 84 POGO-C 846 865 20 15 84 POGO-C 888 899 11 20 239 POGO-C 486 1137 11 20 239 POGO-C 1126 1137 11 20 267 POGO-C 1340 1350 10 15 118 POGO-C 1340 125 15 118 POGO-C 1340 120 15 118 POGO-C 617 632 15 15 158 POGO-C 782 797 15 20 231 POGO-C 782 797 15 20 231 POGO-C 7845 865 20 15 158 POGO-C 7845 865 20 15 164 POGO-C 1340 1350 10 15 178 POGO-C 1340 1350 10 15 188 POGO-C 1340 1350 10 15 188 POGO-C 1340 1350 10 15 158 HULA-A 30 95 5 15 160 HULA-A 174 190 16 20 167 HULA-A 275 290 15 20 174 HULA-A 310 330 20 15 158 HULA-A 359 373 14 20 40 HULA-A 450 465 15 15 150 HULA-A 480 490 10 15 202 HULA-A 480 490 10 15 202 HULA-A 480 490 10 15 205 HULA-A 174 190 16 20 174 HULA-A 480 490 10 15 205 HULA-A 174 190 16 20 191 HULA-A 480 490 10 15 202 HULA-A 174 190 16 20 191 HULA-A 480 490 10 15 202 HULA-A 190 16 20 157 HULA-A 192 125 15 15 100 HULA-A 192 125 15 15 100 HULA-A 192 125 15 20 205 HULA-A 192 125 15 15 131 HULA-B 190 200 10 15 205 15 205 HULA-A 192 20 20 10 15 205 15 205 HULA-A 1078 1083 5 15 205 HULA-A 1
163 HULA-B 190 200 10 15 165 HULA-B 220 270 50 15 27 HULA-B 293 309 16 20

284 HULA-B 1335 1345 10 15 296 HULA-B 1410 1445 35 15
284 HULA-B 1335 1345 10 15 296 HULA-B 1410 1445 35 15 1 COUK-A 11 25 14 20 149 COUK-A 11 25 14 20 149 COUK-A 226 230 5 15 168 COUK-A 226 230 5 15 168 COUK-A 375 405 30 15 193 COUK-A 375 405 30 15 193 COUK-A 455 465 10 15 196 COUK-A 453 560 25 15 205 COUK-A 576 595 20 15 222 COUK-A 750 755 5 15 233 COUK-A 850 880 30 15 237 COUK-A 895 935 40 15 246 COUK-A 1170 1185 15 15 119 COUK-A 1223 1239 16 20 130 COUK-A 1320 1365 45 20 138 COUK-A 1421 1436 15 15 209 COUK-B 315 329 14 20 134 COUK-B 315 329 14 20 134 COUK-B 315 329 14 20 144 COUK-B 315 329 14 20 157 INDI-A 20 25 15 211 COUK-B 715 900 185 15 211 COUK-B 1390 1645 255 15 145 INDI-A 20 25 15 145 INDI-A 20 25 15 145 INDI-A 210 230 20 15 170 INDI-A 138 152 14 20 164 INDI-A 210 230 20 15 171 INDI-A 138 152 14 20 164 INDI-A 210 230 20 15 175 INDI-A 350 400 50 15 180 INDI-A 455 640 5 15 190 INDI-A 456 535 20 15 175 INDI-A 350 50 20 15 175 INDI-A 138 152 14 20 164 INDI-A 200 25 15 175 INDI-A 350 400 50 15 180 INDI-A 456 535 20 15 175 INDI-A 350 50 10 15 180 INDI-A 456 535 20 15 175 INDI-A 350 50 10 15 207 INDI-A 615 620 5 15 207 INDI-A 138 152 14 20 152 INDI-A 128 125 20 15 228 INDI-A 128 129 10 10 20 15 228 INDI-A 133 1318 5 15 229 INDI-A 133 1318 5 15 228 INDI-A 133 1318 5 15 238 INDI-A 133 1318 5 15 238 INDI-A 1337 1425 50 15 249 I
166 CUUR-A 225 230 5 15 168 CDOK-A 245 265 20 15 34 COOK-A 327 343 16 20
183 COOK-A 375 405 30 15 193 COOK-A 455 465 10 15
203 COOK-A 535 560 25 15 205 COOK-A 575 595 20 15
222 COOK-A 750 755 5 15 233 COOK-A 850 880 30 15 237 COOK-A 850 880 30 15
246 COOK-A 975 1020 45 15 259 COOK-A 1095 1110 15 15
268 COOK-A 1170 1185 15 15 119 COOK-A 1223 1239 16 20 130 COOK-A 1320 1365 45 20
138 COOK-A 1392 1406 14 20 293 COOK-A 1421 1436 15 15
2 CUOK-B 17 32 15 20 31 COOK-B 315 329 14 20 184 COOK-B 380 395 15 15
199 COOK-B 510 530 20 15 221 COOK-B 715 900 185 15
91 CUUK-B 933 948 15 20 274 CUUK-B 1225 1270 45 15 291 CUUK-B 1390 1645 255 15
145 INDI-A 20 25 5 15 4 INDI-A 52 66 14 20
12 INDI-A 138 152 14 20 164 INDI-A 210 230 20 15
171 INDI-A 265 285 20 15 170 INDI-A 300 305 5 15 175 INDI-A 320 330 10 15
180 ÎNDI-A 350 400 50 15 185 ÎNDI-A 415 430 15 15
190 INDI-A 450 470 20 15 201 INDI-A 515 535 20 15 204 INDI-A 550 560 10 15
207 INDI-A 575 585 10 15 214 INDI-A 615 620 5 15
219 INDI-A 690 710 20 15 223 INDI-A 760 795 35 15
227 INDI-A 810 850 40 15 228 INDI-A 865 870 5 15 234 INDI-A 885 905 20 15
242 INDI-A 935 965 20 15 245 INDI-A 975 980 5 15
249 INDI-A 995 1015 20 15 253 INDI-A 1040 1050 10 15 256 INDI-A 1080 1100 20 15
262 INDI-A 1120 1145 25 15 269 INDI-A 1170 1215 45 15
269 INDI-A 1120 1145 25 15 269 INDI-A 1170 1215 45 15 273 INDI-A 1230 1240 10 15 123 INDI-A 1282 1298 16 20 282 INDI-A 1313 1318 5 15 280 INDI-A 1333 1338 5 15 288 INDI-A 1375 1425 50 15
280 INDI-A 1333 1338 5 15 288 INDI-A 1375 1425 50 15 295 INDI-A 1440 1460 20 15
$\begin{array}{cccccccccccccccccccccccccccccccccccc$
152 $\hat{B}0SS - A$ 49 113 64 15 11 $BOSS - A$ 135 150 15 20 26 $BOSS - A$ 290 304 14 20 172 $BOSS - A$ 319 339 20 15 189 $BOSS - A$ 405 450 45 15 46 $BOSS - A$ 481 495 14 20 54 $BOSS - A$ 481 495 14 20 54 $BOSS - A$ 580 591 11 20
189 BOSS-A 405 450 45 15 46 BOSS-A 481 495 14 20 54 BOSS-A 529 544 15 20 56 BOSS-A 580 591 11 20
65 BOSS-A 628 643 15 20 216 BOSS-A 660 670 10 15
226 BOSS-A 800 810 10 15 88 BOSS-A 920 931 11 20 238 BOSS-A 946 956 10 15
252 BUSS-A 1019 1036 17 20 252 BUSS-A 1051 1071 20 15 108 BUSS-A 1121 1135 14 20 265 BUSS-A 1150 1155 5 15
252 BOSS-A 1051 1071 20 15 108 BOSS-A 1121 1135 14 20 265 BOSS-A 1121 1135 14 20 213 BOSS-A 1150 1155 5 15 113 BOSS-A 1191 1205 14 20 275 BOSS-A 1240 1250 10 15
$\begin{array}{cccccccccccccccccccccccccccccccccccc$
283 BOSS-A 1336 1346 10 15

Day 7

143 POGO-A	15 30 15 15
153 POGO-A	60 75 15 15
155 POGO-A	90 105 15 15
161 POGO-A	125 145 20 15
164 POGO-A	190 200 10 15
168 POGO-A	245 275 30 15
177 Pogo-A	325 335 10 15
183 POGO-A 191 POGO-A	375 390 15 15
185 POGO-A	450 460 10 15

196 P0G0-A 500 545 45 15 51 P0G0-A 572 588 16 20 208 P0G0-A 603 618 15 15 60 P0G0-A 648 653 15 20
213 POGO-A 678 688 10 15 72 POGO-A 721 732 11 20 227 POGO-A 780 785 5 15 231 POGO-A 835 850 15 15 85 POGO-A 878 895 17 20
88 POGO-A 921 931 10 20 242 POGO-A 946 961 15 15 98 POGO-A 992 1004 12 20 251 POGO-A 1020 1035 15 15 254 POGO-A 1050 1060 10 15 107 POGO-A 1050 1060 10 15
107 POGO-A 1091 1104 13 20 260 POGO-A 1119 1129 10 15 268 POGO-A 1155 1165 10 15 116 POGO-A 1155 1212 17 20 119 POGO-A 1249 1264 15 20 126 PUGO-A 1297 1313 16 20 282 POGO-A 1368 1384 16 20 284 POGO-A 1368 1384 16 20
1282 POGO-A 1330 1342 12 15 134 POGO-A 1368 1384 16 20 295 POGO-A 1410 1420 10 15 291 POGO-A 1435 1445 10 15 147 POGO-B 35 70 35 15 158 POGO-B 115 130 15 15
168 POGO-B 115 130 15 15 160 POGO-B 145 160 15 15 163 POGO-B 145 160 15 15 170 POGO-B 265 275 10 15 170 POGO-B 265 275 10 15 179 POGO-B 330 375 45 15
190 POGO-B 405 450 45 15 193 POGO-B 465 510 45 15 47 POGO-B 550 561 11 20 52 POGO-B 590 606 16 20 206 POGO-B 621 636 15 15
66 P0G0-B 674 690 16 20 218 P0G0-B 705 720 15 15 73 P0G0-B 745 759 14 20 80 P0G0-B 821 831 10 20 82 P0G0-B 851 867 16 20 235 P0G0-B 895 910 15 15
22 POGO-B 938 952 14 20 99 POGO-B 994 1010 16 20 247 POGO-B 1025 1035 10 15 106 POGO-B 1081 1097 16 20 110 POGO-B 1118 1131 13 20 112 POGO-B 1159 1172 13 20
206 POGO-B 621 636 15 15 66 POGO-B 674 690 16 20 218 POGO-B 705 720 15 15 73 POGO-B 745 759 14 20 80 POGO-B 821 831 10 20 82 POGO-B 851 867 16 20 235 POGO-B 895 910 15 15 92 POGO-B 994 1010 16 20 247 POGO-B 1025 1035 10 15 106 POGO-B 1081 1097 16 20 110 POGO-B 1159 1172 13 20 215 POGO-B 1159 1172 13 20 215 POGO-B 1289 1304 15 20 132 POGO-B 1289 1304 15 20 132 POGO-B 1348 1363 15 20 140 POGO-B 1400 1415 15 20
140 POGD-B 1400 1415 15 20 157 POGD-C 90 135 45 15 162 POGD-C 240 255 15 15 173 POGD-C 285 330 45 15 184 POGD-C 375 405 30 15 194 POGD-C 375 405 30 15 194 POGD-C 480 490 10 15 200 POGD-C 530 555 25 15 53 POGD-C 658 614 16 20 62 POGD-C 651 663 12 20 68 POGD-C 691 707 16 20 74 POGD-C 751 766 15 20 232 POGD-C 840 860 20 15 234 POGD-C 840 860 80 10 15 24 POGD-C 840 860 80 10 15 25 POGD-C 840 860 80 10 15 25 POGD-C 840 860 80 15 25 POGD-C 840 860 80 10 15 25 POGD-C 840 80 10 15 25 POGD-C 840 80 15 25 POGD-C 840 80 15 25 POGD-C 840 80 15 25 POGD-C 840 80 15 25 POGD-
162 POGO-C 165 180 15 15 167 POGO-C 240 255 15 15 173 POGO-C 285 330 45 15 184 POGO-C 375 405 30 15 194 POGO-C 480 490 10 15 200 POGO-C 530 555 25 15 53 POGO-C 598 614 16 20 62 POGO-C 651 663 12 20 68 POGO-C 691 707 16 20 74 POGO-C 751 766 15 20 232 POGO-C 840 860 20 15
244 POGO-C 975 1020 45 15 108 POGO-C 1094 1111 17 20
114 POGO-C 1182 1196 14 20 118 POGO-C 1230 1244 14 20 121 POGO-C 1269 1284 15 20
149 HULA-A 90 100 10 15 156 HULA-A 90 100 10 15
12 HULA-A 101 177 10 20 165 HULA-A 210 215 5 15 22 HULA-A 262 277 15 20 171 HULA-A 262 297 5 15
27 HULA-A 317 329 12 20 32 HULA-A 359 374 15 20 181 HULA-A 389 404 15 15 38 HULA-A 443 455 12 20

262 INDI-A 1115 1190 75 15 274 INDI-A 1210 1225 15 15 120 INDI-A 1252 1268 16 20 279 INDI-A 1283 1293 10 15 287 INDI-A 1350 1370 20 15 286 INDI-A 1385 1395 10 15 286 INDI-A 1415 1435 20 15	
$ \begin{array}{c} 1 & 0.5S-A & 30 & 146 & 16 & 20 \\ 3 & B0SS-A & 130 & 145 & 15 & 20 \\ 21 & B0SS-A & 260 & 273 & 13 & 20 \\ 172 & B0SS-A & 288 & 293 & 5 & 15 \\ 176 & B0SS-A & 320 & 330 & 10 & 15 \\ 31 & B0SS-A & 359 & 375 & 16 & 20 \\ 187 & B0SS-A & 390 & 405 & 15 & 15 \\ 189 & B0SS-A & 390 & 405 & 5 & 15 \\ 189 & B0SS-A & 20 & 425 & 5 & 15 \\ \end{array} $	
189 BOSS-A 420 425 5 15 40 BOSS-A 463 475 12 20 44 BOSS-A 502 512 10 20 46 BOSS-A 539 614 15 20 54 BOSS-A 539 614 15 20 210 BOSS-A 630 640 10 15 64 BOSS-A 667 681 14 20 219 BOSS-A 696 701 5 15 221 BOSS-A 696 701 0 15	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	
151 BUSS-B 60 360 300 15 45 BUSS-B 512 527 15 20 49 BUSS-B 560 574 14 20 55 BUSS-B 600 617 17 20 205 BUSS-B 632 652 20 15 69 BUSS-B 700 716 16 20 97 BUSS-B 989 1005 16 20	
130 BOSS-B 1323 1336 13 20 137 BOSS-B 1380 1395 15 20 4 LION-A 59 75 16 20 7 LION-A 124 137 13 20 13 LION-A 167 178 11 20 19 LION-A 220 235 15 20 166 LION-A 250 270 20 15 26 LION-A 294 309 15 20 175 LION-A 294 329 5 15	
26 LION-A 294 309 15 20 175 LION-A 324 329 5 15 33 LION-A 361 374 13 20 186 LION-A 390 395 5 15 198 LION-A 515 535 20 15 65 LION-A 670 679 9 20 217 LION-A 694 714 20 15 77 LION-A 786 802 16 20 81 LION-A 894 845 11 20 90 LION-A 929 943 14 20	
217 LION-A 786 802 16 20 81 LION-A 786 802 16 20 81 LION-A 786 802 16 20 81 LION-A 929 943 14 20 238 LION-A 929 943 14 20 238 LION-A 929 943 14 20 238 LION-A 929 943 14 20 257 LION-A 1027 1040 13 20 257 LION-A 1080 1100 20 15 272 LION-A 1207 1217 10 15 123 LION-A 1285 1297 12 20 131 LION-A 1334 1349 15 20 135 LION-A 1372 1382 10 20 297 LION-A 1423 1433 10 15 146 TON-PA 30 53 21 5	
146 LION-B 30 53 23 15 23 LION-B 263 278 15 20 28 LION-B 321 331 10 20 34 LION-B 395 412 17 20 42 LION-B 497 510 13 20 67 LION-B 685 700 15 20 226 LION-B 780 840 60 15 89 LION-B 929 943 14 20 246 LION-B 990 1010 20 15	
20 GUAN-A 250 305 55 20 30 GUAN-A 350 364 14 20	
37 GUAN-A 428 441 13 20 41 GUAN-A 466 478 12 20 50 GUAN-A 565 581 16 20 58 GUAN-A 615 629 14 20 83 GUAN-A 655 869 14 20	

Bibliography

- 1. Arbabi, Mansur and John A. Garate. "Interactive Real Time Scheduling and Control." Proceedings of the 1985 Summer Simulation Conference. 271-277. 1985.
- 2. Baker, Bruce N. Introduction to Sequencing and Scheduling. John Wiley & Sons, 1974.
- Blanton Jr., Joe L. and Roger L. Wainwright. "Multiple Vehicle Routing with Time and Capacity Constraints using Genetic Algorithms." Proceedings of 5th International Conference on Genetic Algorithms, edited by S. Forrest. San Mateo, CA: Morgan Kauffman, 1993.
- 4. Bruns, Ralf. "Direct Chromosome Representation and Advanced Genetic Operators for Production Scheduling." Proceedings of 5th International Conference on Genetic Algorithms, edited by S. Forrest. San Mateo, CA: Morgan Kauffman, 1993.
- 5. Davis, Lawerence. The Handbook of Genetic Algorithms. New York: Van Nostrand Reinhold, 1991.
- 6. Garey, Michael R. and David S. Johnson. Computers and Intractability: A Guide to the Theory of NP-Completeness. San Francisco: W.H. Freeman, 1979.
- 7. Goldberg, David. Genetic Algorithms in Search, Optimization, and Machine Learning. Reading, Mass.: Addison-Wesley, 1989.
- 8. Gooley, Capt Timothy. Automating the Satellite Range Scheduling Process. MS thesis, AFIT/GOR/ENS/93M-06, School of Engineering, Air Force Institute of Technology (AU), Wright-Patterson AFB OH, March 1993. (AAK-1926).
- 9. Grefenstette, John, Rajeev Gopal, Brian Rosmaita, and Dirk Van Gucht. "Genetic Algorithms for the Traveling Salesman Problem." Proceedings of an International Conference on Genetic Algorithms and their Applications, edited by John Grefenstette. Hillsdale, NJ: Lawrence Erlbaum Associates, 1985.
- 10. Holland, John. Adaptation in Natural and Artificial Systems. Ann Arbor: University of Michigan Press, 1975. reprinted by MIT Press, Cambridge, MA, 1992.
- Kennedy, Capt Dale J. A Prototype Expert System Advisor for Satellite Support Scheduling. MS thesis, AFIT/GOR/OS/86D-5, School of Engineering, Air Force Institute of Technology (AU), Wright-Patterson AFB OH, March 1986. (AD-A179 425).
- 12. Lawler, E.L., editor. The Traveling Salesman Problem. UK: John Wiley and Sons, 1985.
- 13. Michalewicz, Zbigniew. Genetic Algorithms + Data Structure = Evolution Programs. New York: Springer-Verlag, 1992.
- Schalck, Capt Stanley Michael. Automating Satellite Range Scheduling. MS thesis, AFIT/GSO/ENS/93D-14, School of Engineering, Air Force Institute of Technology (AU), Wright-Patterson AFB OH, December 1993. (AD-A273 829).

- 15. Starkweather, T. and Darryl Whitley. "A Comparison of Genetic Sequencing Operators." Proceedings of 4th International Conference on Genetic Algorithms, edited by R.K. Belew and L.B. Booker. Los Altos, CA: Morgan Kaufmann, 1991.
- 16. Syswerda, Gilbert. The Handbook of Genetic Algorithms, chapter 21. New York: Van Nostrand Reinhold, 1991. Schedule Optimization Using Genetic Algorithms.
- 17. Whitley, Darrell. "GENITOR: a different genetic algorithm." Proceedings of the Rocky Mountain Conference on Artificial Intelligence. 118-130. 1988.
- 18. Whitley, Darrell, Timothy Starkweather, and Daniel Shaner. The Handbook of Genetic Algorithms, chapter 22. New York: Van Nostrand Reinhold, 1991. The Traveling Salesman and Sequence Scheduling: Quality Solutions using Genetic Edge Recombination.
- 19. Whitley, Darryl, Timothy Starkweather, and D'Ann Fuquay. "Scheduling Problems and Traveling Salesman: The Genetic Edge Recombination Operator." Proceedings of 3rd International Conference on Genetic Algorithms, edited by J.D. Schaffer. Los Altos, CA: Morgan Kaufmann, 1989.

REPORT DOCUMENTATION PAGE			Form Approved OMB No. 0704-0188
athering and maintaining the data needed ollection of information in huding suggestic	and completing and reviewing the collection	of information. Send comments rega	Purewing instructions searching existing data source inding this burden estimate or any other aspect of thi information Operations and Reports (215) jefferso ject (0704-0188), Washington (2) 20503
I. AGENCY USE ONLY (Leave b)		3. REPORT TYPE AN Master's These	
4. TITLE AND SUBTITLE A GENETIC ALGORITHM APPROACH TO AUTOMATING SATELLITE RANGE SCHEDULING		5. FUNDING NUMBERS	
i. AUTHOR(S) Donald A. Parish, Capt,	USAF		
7. PERFORMING ORGANIZATION NAME(S) AND ADDRESS(ES) Air Force Institute of Technology, WPAFB OH 45433-6583		8. PERFORMING ORGANIZATION REPORT NUMBER AFIT/GOR/ENS/94M-10	
9. SPONSORING / MONITORING AGENCY NAME(S) AND ADDRESS(ES)		10. SPONSORING / MONITORING AGENCY REPORT NUMBER	
1. SUPPLEMENTARY NOTES			
2a. DISTRIBUTION / AVAILABILITY Distribution Unlimited	Y STATEMENT		12b. DISTRIBUTION CODE
 Distribution Unlimited ABSTRACT (Maximum 200 wo Satellite range schedulin tracking station commu supports continue to ince efficiently. Previous rese may not produce the be approach to automating was to schedule as many to find the best priority schedules based on simp data with slightly better Based on the reported 	rds) ng involves scheduling satell nicate with each other withis crease, more pressure is place earch efforts focused on heur est results. The objective of the generation of 24 hour so y supports as possible witho y ordering of support reques ple rules. A schedule was put results compared to earlier	in a specified time windo ed on the current manua stic and mixed-integer p f this research was to de thedules was competitive ut conflict. The genetic a ts, and then used a sche roduced for seven days of results using a mixed-inter n approach presented in	12b. DISTRIBUTION CODE satellite and a specific remote w. As the number of satellite l system to generate schedules rogramming approaches which termine if a genetic algorithm with other methods. The goal algorithm approach attempted dule builder program to build f representative satellite range eger programming formulation. this research appears to be a
 Distribution Unlimited ABSTRACT (Maximum 200 wo Satellite range schedulin tracking station commu supports continue to ince efficiently. Previous rese may not produce the be approach to automating was to schedule as many to find the best priority schedules based on simp data with slightly better Based on the reported competitive approach for SUBJECT TERMS 	rds) ng involves scheduling satell nicate with each other withit crease, more pressure is place earch efforts focused on heuri est results. The objective of the generation of 24 hour so y supports as possible witho y ordering of support reques ple rules. A schedule was pr results compared to earlier results, the genetic algorithm	in a specified time windo ed on the current manua istic and mixed-integer p f this research was to de chedules was competitive ut conflict. The genetic a ts, and then used a sche roduced for seven days of results using a mixed-inte m approach presented in e range schedules.	satellite and a specific remote w. As the number of satellite l system to generate schedules rogramming approaches which termine if a genetic algorithm with other methods. The goal algorithm approach attempted dule builder program to build f representative satellite range eger programming formulation.

• ` ٠

Prescribed by ANSI Std. 239-18 298-102