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CITATION:

Tsuboi, Akito. Colored Random Graphs and the Order Property (Model theoretic aspects of the notion of independence and dimension). 数理解析研究所講究録 2023, 2249: 64-67

ISSUE DATE:

2023-04

URL:

<http://hdl.handle.net/2433/285439>

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# Colored Random Graphs and the Order Property

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## 1 Introduction

In this article, a graph means an  $R$ -structure, where  $R$  is a binary symmetric irreflexive predicate. If  $R(a, b)$  holds, we consider  $a$  and  $b$  are adjacent by an edge. A subgraph means a substructure, in the graph theory terminology, it is an induced subgraph. A finite coloring of a graph  $G$  usually means a function  $f : R^G \rightarrow F$ , where  $F$  is a finite set of colors. However, we are going to take a slightly different setting, which will be explained later. A monochromatic subgraph is a subgraph  $H$  for which the coloring function  $f$  is constant on  $R^H$ . In general, it is an important question whether a colored graph has monochrome subgraphs of a certain kind. Here we concentrate on countable random graphs and their coloring.

A graph  $G$  is called a random graph, it satisfies the following axioms for all disjoint subsets  $A \neq \emptyset$  and  $B$ ,

$$\exists x \left( \bigwedge_{a \in A} R(a, x) \wedge \bigwedge_{b \in B} \neg R(a, b) \right).$$

A random graph is necessarily infinite, and is universal in the sense that it embeds all finite graphs. It is easy to see that the theory of a random graph is  $\aleph_0$ -categorical, and is simple. In [2], they proved:

- (\*) A colored countable random graph  $G$  has a subgraph  $H$  such that  $H \cong G$  (as graphs) and that  $H$  is (at most) 2-colored.

They also gave an example of  $G$  without monochromatic subgraph  $H \cong G$ .

In this article, we study the case when  $G$  does not have a monochromatic subgraph  $H \cong G$ . As a main result, we state some relation between the

coloring and the instability strength (see Theorem 9). We do not give a detail of the proof.

## 2 Definitions and Preliminaries

Let  $G$  be a countable random graph in the language  $\{R\}$ , where  $R$  is a binary predicate symbol for edges.

Let  $N \in \omega$ . An  $N$ -coloring of  $G$  means an expansion of  $G$  to the language  $L \cup \{R_i\}_{i < N}$  such that  $R^G$  is the disjoint union of  $R_i^G$  ( $i < N$ ). For a subset  $C \subset N$ ,  $R_C(x, y)$  is an abbreviation of  $\bigvee_{i \in C} R_i(x, y)$ . If  $R_i(a, b)$  holds, we think that the edge  $ab$  is painted in the color  $i$ .

Now we fix a countable random graph  $G$ . We assume the edges of  $G$  are  $N$ -colored.

$S_{na}$  denotes the set of all non-algebraic types with a finite domain.

**Definition 1.** 1. Let  $p \in S_{na}$ . An infinite subset  $X \subset G$  is  $p$ -large, if (1)  $p(X) = X$  and (2)  $q(X)$  is non-empty for all non-algebraic  $q \supset p$ . We say  $X$  is large, if  $p(X)$  is  $p$ -large for some  $p$ .

2. We write  $X \subset_{\text{lrg}} Y$ , if  $X \subset Y$  and  $X$  is large.

Then, we can prove the following lemmas. (Proofs are not shown here.)

**Lemma 2.** *Suppose that  $X$  is  $p$ -large and that  $X = \bigcup_{i < n} X_i$ , where  $n \in \omega$ . Then, there is an index  $i < n$  and a non-algebraic type  $q \supset p$  such that  $q(X_i)$  is  $q$ -large.*

**Lemma 3.** *Suppose that  $X$  and  $Y$  are large. Then, there is a color  $i < N$  and  $X_0 \subset_{\text{lrg}} X$  such that, for all  $a \in X_0$ , both*

$$\{b \in Y : R_i(a, b)\} \text{ and } \{b \in Y : \neg R(a, b)\}$$

*are large.*

**Definition 4.** Let  $X$  and  $Y$  be large.

1.  $C(X, Y)$  denotes the set of all colors  $i < N$  for which some  $X_0 \subset_{\text{lrg}} X$  satisfies the statement of Lemma 3.
2.  $C^*(X, Y) = \bigcap \{C(X', Y') : X' \subset_{\text{lrg}} X, Y' \subset_{\text{lrg}} Y\}$ .

**Lemma 5.** *Let  $X, Y$  be large. Then, there is  $X_0 \subset_{lr} X$  and  $Y_0 \subset_{lr} Y$  such that  $C^*(X_0, Y_0) \neq \emptyset$ .*

**Lemma 6.** *There is a large set  $Z$  and  $i^*, j^* < N$  such that for any large  $W \subset Z$  there is a disjoint large sets  $X, Y \subset W$  such that  $i^* \in C^*(X, Y)$  and  $j^* \in C^*(Y, X)$ .*

### 3 Main Results

Now we fix a large set  $Z$  and  $i^*, j^* < N$  satisfying the requirement in Lemma 6.

**Definition 7.** Let  $A$  be a finite subset of  $Z$ , and  $D \supset A$  a finite subset of  $G$ . Let  $\mathfrak{X} = \{X_p\}_{p \in S_{na}(A)}$  be a set of large subsets of  $Z$  and let  $\mathfrak{T} = \{p^*\}_{p \in S_{na}(A)}$  be a set of types. We say the tuple  $(A, D, \mathfrak{X}, \mathfrak{T})$  is good, if the following are true: For all  $p \neq q \in S_{na}(A)$ , 1.  $p \subset p^* \in S_{na}(D)$ ; 2.  $X_p$  is  $p^*$ -large; 3.  $(i^*, j^*)$  or  $(j^*, i^*)$  belongs to  $C^*(X_p, X_q) \times C^*(X_q, X_p)$ . 4. For all  $a \in A$  and  $b \in AX_p$ ,  $R(a, b) \iff R_{\{i^*, j^*\}}(a, b)$ .

**Proposition 8.** *Suppose that  $(A, D, \{X_p\}_{p \in S_{na}(A)}, \{p^*\}_{p \in S_{na}(A)})$  is good. Then, for all  $s \in S_{na}(A)$ , we can find  $d \in X_s$ ,  $D' \supset D$ ,  $\{X_q\}_{q \in S_{na}(Ad)}$  and  $\{q^*\}_{q \in S_{na}(Ad)}$  such that*

- $(Ad, D', \{X_q\}_{q \in S_{na}(Ad)}, \{q^*\}_{q \in S_{na}(Ad)})$  is also good;
- $p^* \subset q^*$  and  $X_q \subset X_p$ , if  $p \in S_{na}(A)$ ,  $q \in S_{na}(Ad)$  and  $p \subset q$ .

**Theorem 9.** *Let  $G$  be a random graph and suppose that an  $N$ -coloring is given on  $G$  by  $L^* = \{R, R_1, \dots, R_N\}$ . Then the following conditions are equivalent:*

- (a)  $G$  does not have a monochromatic generic subgraph;
- (b) For any generic subgraph  $G_0 \subset G$ , there is a generic  $H \subset G_0$  having the strict order property in the expanded language  $L^*$ .

*Sketch of Proof.* (b)  $\Rightarrow$  (a): This is trivial since a monochromatic subgraph is a mere random graph. (a)  $\Rightarrow$  (b): We assume (a). For simplicity of the notation, we can assume  $G_0 = G$ . We choose  $i^*, j^* < N$  and  $Z$  as in Lemma 6. Since  $G$  does not have a monochromatic generic subgraph, we have  $i^* \neq j^*$ . So, for simplicity, we assume  $i^* = 0$  and  $j^* = 1$ . Let  $\{g_i\}_{i \in \omega}$  be an enumeration of  $G$  such that for all  $i > 0$ ,

1.  $R(g_0, g_i)$  if and only if  $i$  is even;
2.  $R(g_{4i}, g_j)$  for all odd numbers  $j < 4i$ .

Notice that such an enumeration does exist. Choose disjoint large subsets  $X_0, X_1 \subset Z$  such that  $0 \in C^*(X_0, X_1)$  and  $1 \in C^*(X_1, X_0)$ . We are going to define  $h_i$  ( $i < \omega$ ) such that  $(g_i)_{i \in \omega} \cong (h_i)_{i \in \omega}$ . By symmetry, shrinking  $X_0$  and  $X_1$ , we may assume  $\forall x \in X_0 (R(h_0, x))$  and  $\forall x \in X_1 (\neg R(h_0, x))$  hold for some  $h_0 \in G$ . In this proof, we inductively choose elements  $h_i \in X_0 X_1$  ( $i > 0$ ) such that, by letting

$$D_n := \{h_m : h_m \neq h_0, \neg R(h_0, h_m) \text{ and } R_1(h_{4n}, h_m)\},$$

$\{D_n : n \in \omega\}$  forms a strictly increasing sequence of uniformly defined sets. Thus,  $H := \{h_i\}_{i \in \omega}$  has the strict order property.  $\square$

## References

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