

Building a Secure Image Cryptography System using Parallel Processing and Complicated Dynamic Length Private Key

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Abstract: A method of color images cryptography will be introduced, programmed, and tested. The proposed method is based on using a digital color image as an image key; this image is to be kept secret without transmission. The proposed method will provide a high level of images protection based on the complicated and complex private key used in cryptography, this key will be changed when replacing the image key, or changing the data block size, or changing the color channel. The proposed method will be compared with other standard methods of data cryptography, and it will be shown how this method will improve the efficiency of data cryptography by minimizing the encryption-decryption time, the obtained results will be compared with the standard method of data cryptography to show the speedup achieved by the proposed method. It will be shown how to execute the proposed method in parallel, 2, 4, and 8 threads will be used to execute the method and the associated speedup will be calculated. The proposed method will protect the data by providing a high level of security, this can be achieved by using a variable-length private key, the private key length and content will depend on the selected image key, selected color matrix, and the selected block size. The block size used in the proposed method will be variable and it will be shown that the proposed method will satisfy the quality requirements by providing good value for Mean Square Error (MSE), and Peak Signal to Noise Ratio (PSNR)

Keywords: Image_key; PK; block size; MSE; PSNR

1 Introduction

Color images are considered one of the most widespread digital data types used in many critical vital applications. This spread is due to several reasons, the most important of which are [1,2,3,4]:

- Ease of obtaining a digital color image at no cost.
- The high image size can be used for different purposes and take advantage of the digital data contained in the digital image.
- The text in the entries may be of any length.
- The ease of processing a color digital image is usually represented by a three-dimensional matrix, one for each of the three colors (red, green, and blue) and as shown in Figure 1.

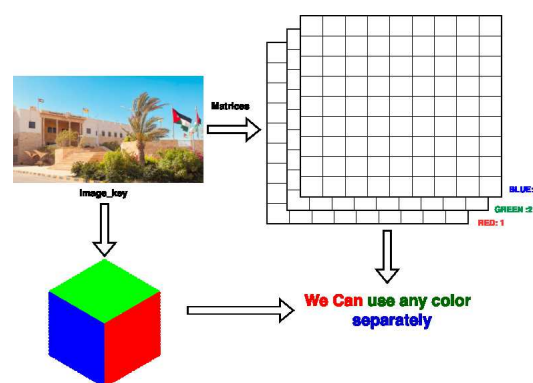


Fig. 1: Color image representation

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- The possibility of dealing with each color's matrix and parts of the image separately.
- The possibility of adjusting the size of the image to suit any other size, whether by reducing the size or increasing it, is shown in Figure 2.

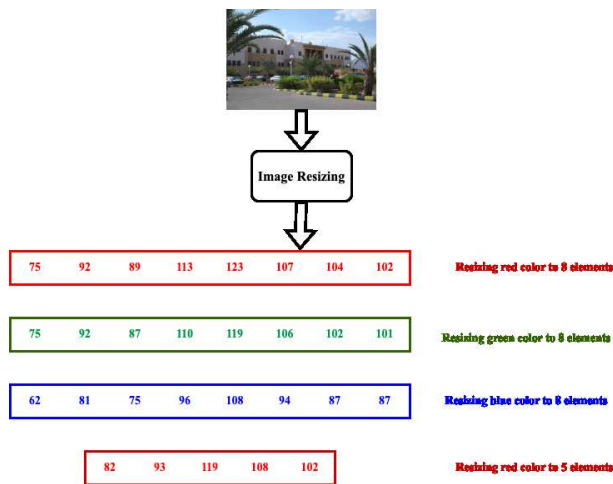


Fig. 2: Image resizing

The possibility of using digital images to protect all types of digital data using hiding (data steganography), encryption, and decryption processes (data cryptography). Color digital images require protection from the danger of tampering, intruders, or data thieves for several reasons [5,6,7,8], the most important of these reasons. First, the digital image can be of a private or confidential nature. Second is the possibility of the digital image carrying personal data. Third, frequent use of digital images in applications requires more protection. Fourth, the circulation of digital images might provide the possibility of accessing them by unrelated or unauthorized persons or entities. Data cryptography is one method used to protect digital color images; it can be performed as shown in Figure 3 by using the private key (PK) and manipulating several arithmetical and logical operations using the data to be encrypted and PK.

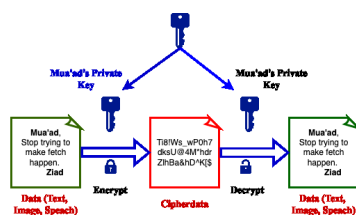


Fig. 3: Data cryptography process

2 Related Work

Multiple methods [9,10,11,12,13,14] encrypt and decrypt data, including digital images. When choosing a specific method to protect data, this method must achieve the following conditions:

- The method should be secure to provide a high degree of data protection, achieved using a private secret key that is difficult to hack, know or guess.
- The method works to destroy and distort the data when encrypting so that the data becomes useless or difficult to understand for any unauthorized third party and works to return the original data without change when decrypting. The percentage of destruction can be measured using the quality parameter MSE or PSNR; the MSE value between the encrypted data and the original one must be very high (PSNR must be very low, while MSE between the original data and the decrypted one must equal zero (PSNR must be infinite) (see equations 1, and 2) [9,10].

$$MSE = \frac{1}{mn} \sum_{i=1}^m \sum_{j=1}^n (x_{ij} - y_{ij})^2 \quad (1)$$

where m is the number of rows in cover image, n is the number of columns in cover image, x_{ij} is the pixel value from cover image, and y_{ij} is the pixel value from stego image [10,11,12].

$$PSNR = 10 \log_{10} \frac{[MAX_I]^2}{MSE_t} \quad (2)$$

where MAX_I is the maximum signal value that exists in our original "known to be good" image [9].

- MSE and PSNR are good parameters to measure the changes between the source image and the encrypted/decrypted image, the changes are due to applying encryption-decryption process, If MSE equal zero, then PSNR is equal infinite, this means that there is no changes in the decrypted image, and the decrypted image is identical to the source image [10].
- The method should be flexible, so it is easy to change the private key or modify it by increasing or decreasing its length. It is also easy to alter the data block length used in encryption and decryption [11].
- The method should be highly efficient by reducing the encryption time and decryption times to the lowest possible value, thus increasing the method throughput (number of bytes treated in a unit of time) [11].
- The method should be easy and feasible, programmatically or hardware [12].

Multiple methods are now used to protect data based on international standards, including Data Encryption standard (DES), Triple DES (3DES), Advanced Encryption Standard (AES), and blowfish (BF) [15,16,17,18,19]. These methods share many inefficient and

unsecured features; Table 1 shows these methods' main features.

Faster methods were introduced; these methods were used to minimize the encryption-decryption times and to maximize the throughput of data cryptography; in [30], the authors provided a robust and fast image encryption scheme based on a mixing technique. In [31], the authors provided a cosine-transform-based chaotic system for image encryption. In contrast, in [32], the authors introduced a novel image encryption algorithm based on a polynomial combination of chaotic maps and dynamic function generation. In [33], the authors introduced Multiple-image Encryption Algorithm Based on DNA Encoding and Chaotic Systems, while in [34], the authors produced multiple-image encryption with bit-plane decomposition and chaotic maps; these methods provided good quality.

In [21], a comparative analysis of DES, 3DES, and AES was done, and the performance and throughputs of these methods were calculated; the throughput will rapidly decrease when the data size increases; thus, using these methods for image encryption-decryption will be not efficient. In [18, 22, 23, 29, 35], a performance analysis or blowfish method of data cryptography was done; the results showed that using this method will increase the efficiency compared with DES and AES methods, but using it for image encryption-decryption requires enhancements to increase the cryptography process throughput.

3 The Proposed Method

The proposed method shown in Figure 4 provides a high degree of data security and protection through the use of a color image to generate the private keys used in the encryption and decryption process that are difficult to penetrate or guess for the following reasons:

- The key image is determined by agreement between the sender and receiver and is kept secretly and without resorting to sending or circulating it.
- The ability to change the key image at any time and if needed.
- Changing the data block size and the selected color matrix for resizing will change the length and contents of the private key as shown in Figure 5 to Figure 7.
- Changing the image key will change the contents of the private key (see Figure 5 and Figure 6).

The Feistel functions use a rotate operation; the selected number of rotating digits can change from 0 to 7 in the encryption phase and from 7 to 0 in the decryption phase. The proposed method of color image encryption can be implemented, as shown in Figure 8, by applying the following steps:

- Step 1: The initialization step includes the following substeps: select the image key, select the color matrix

to be used to generate PK, select the data block size in bytes, and resize the color matrix to meet the block size; the resized image will be used as a PK, and reshape the image to be encrypted into a one-row matrix.

- Step 2: For each block of data, apply the Feistel function for each byte, XOR the results with the associated byte from the PK, then use the second Feistel function.

- Step 3: Reshape the encrypted data to the 3D matrix to get the encrypted image. The decryption phase can be implemented reversely, as shown in Figure 8.

4 Implementation and Experimental Results

The proposed method was programmed using Matlab code; the program was executed several times using an I7 multicore processor with 8 G byte RAM. Figure 9 shows an output example of the proposed method execution.

Images shown in Table 2 were taken, encrypted decrypted using various initial stages; Table 3 shows the obtained results using image 4 as an image_key, block size=16 bytes:

The same images were taken, encrypted-decrypted using image 2 (small image) as an image_key, with block size=100 bytes; the results are shown in Table 3.

The selected images were encrypted-decrypted, varying the block size. Table 4, Table 5, and Table 6 show the experimental results.

Based on the exciting findings, the encryption time increases as the number of keys increases. The results have shown how the proposed method would improve the efficiency of data cryptography by reducing the encryption-decryption time. The obtained results are compared with the standard method of data cryptography to show the speedup achieved by the proposed method. The traditional cryptography methods that were examined are (DES, 3DES, AES, and BF). The proposed method satisfies the requirements for image quality in the encryption and decryption phases [14, 15].

For comparison purposes, the standard methods of data encryption-decryption were implemented using the same selected images; Table 7 shows the obtained experimental results compared to the methods (DES and 3DES). Also, Table XI compares AES and BF methods and the proposed approach (MPK).

The proposed method was implemented using a multithreading environment by executing the code of the proposed method using Matlab pool as shown in Figure 10, varying the number of threads; the obtained results are shown in Table 8 compared to the studies [24, 25, 26].

The average speedup by implementing 8 threads outperforms the other executed threads. Using one thread is more costly and more time-consuming. Also, increasing the number of threads increases the speed, as shown in Table 8. The multithread approach [27] affects the encryption systems.

Table 1: Cryptography methods main features [1,3,4,5,6,7,36]

| Algorithm parameter | Data Encryption standard (DES) | Triple DES (3DES) | Advanced Encryption Standard (AES) | Blowfish |
|-----------------------------|---------------------------------------|---|---------------------------------------|---------------------------------------|
| PK length (bit) | 56 (fixed) | 112, 168 (fixed) | 128, 192, 256 (fixed) | 32-448 (fixed) |
| Block size (bit) | 64 (fixed) | 64 (fixed) | 128 (fixed) | 64 (fixed) |
| Ability to deal with images | Difficult | Difficult | Difficult | Difficult |
| Encryption quality | Excellent: High MSE and low PSNR | Excellent: High MSE and low PSNR | Excellent: High MSE and low PSNR | Excellent: High MSE and low PSNR |
| Decryption quality | Excellent: Zero MSE and infinite PSNR | Excellent: Zero MSE and infinite PSNR | Excellent: Zero MSE and infinite PSNR | Excellent: Zero MSE and infinite PSNR |
| Efficiency | Slow | Slow | Slow | Moderate |
| Attack | Brute force attack | Brute force attack, Known plaintext, Chosen plaintext | Side channel attack | Dictionary attack |
| Structure | Feistily | Feistel | Substitution-Permutation | Feistel |
| Block cipher | Binary | Binary | Binary | Binary |
| Rounds | 16 (fixed) | 48 (fixed) | 10,12,14 (fixed) | 16 (fixed) |
| Flexibility to modification | no | yes | yes | yes |
| Simplicity | no | no | no | no |
| Security level | Adequate | Adequate | Excellent | Excellent |
| Throughput | Low | low | Low | Moderate |

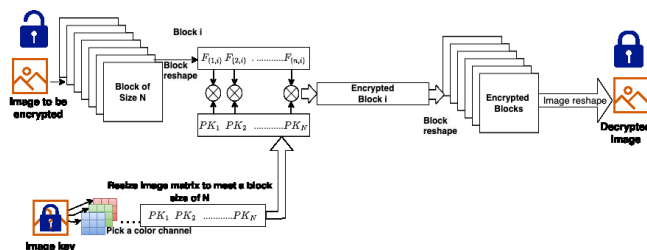


Fig. 4: The proposed method

Table 2: Results using image 4 as an image_key, block size=16 bytes

| Image number | Dimension | Size (byte) | MSE | PSNR | Encryption time (second) |
|--------------|-------------|-------------|-------------|---------|--------------------------|
| 1 | 151x333x3 | 150849 | 1.4272e+004 | 15.1648 | 0.04650 |
| 2 | 152x171x3 | 77976 | 1.5070e+004 | 14.6209 | 0.037891 |
| 3 | 360x480x3 | 518400 | 1.3756e+004 | 15.5326 | 0.094956 |
| 4 | 1071x1600x3 | 5140800 | 1.2969e+004 | 16.1223 | 0.767682 |
| 5 | 981x1470x3 | 4326210 | 1.2915e+004 | 16.1640 | 0.614167 |
| 6 | 165x247x3 | 122265 | 1.2539e+004 | 16.4590 | 0.043193 |
| 7 | 360x480x3 | 518400 | 1.5040e+004 | 14.6404 | 0.098163 |
| 8 | 183x275x3 | 150975 | 1.3846e+004 | 15.4681 | 0.044476 |
| 9 | 183x275x3 | 150975 | 1.2767e+004 | 16.2791 | 0.042990 |
| 10 | 201x251x3 | 151353 | 1.4221e+004 | 15.2005 | 0.042769 |
| 11 | 600x1050x3 | 1890000 | 1.3557e+004 | 15.6790 | 0.299206 |
| 12 | 1144x1783x3 | 6119256 | 1.2623e+004 | 16.3927 | 0.879196 |

5 Discussion

From the obtained experimental results, we can raise many significant findings. First, the proposed method is highly secure; the private key is variable, and the length and the contents of PK depend on the selected image_key, the selected color matrix for resizing, and the block size. Second, the data block size is variable. Third, the

encryption time will decrease when selecting image_key with a small size and using data blocks with bigger sizes, as shown in Figure 11. Fourth, increasing block size will decrease encryption time as in the study [28]; the optimal block size for encrypting and decrypting [29] the selected images was 800 bytes. Fifth, the proposed method provides a significant speedup of the process of data cryptography compared with standard methods. In

Table 3: Results using image 2 as an image_key, block size=100 bytes

| Image Number | Dimension | Size (byte) | MSE | PSNR | Encryption time (second) |
|--------------|-------------|-------------|-------------|---------|--------------------------|
| 1 | 151x333x3 | 150849 | 2.7697e+004 | 8.5344 | 0.009153 |
| 2 | 152x171x3 | 77976 | 3.5373e+004 | 6.0882 | 0.004588 |
| 3 | 360x480x3 | 518400 | 2.1046e+004 | 11.2808 | 0.016747 |
| 4 | 1071x1600x3 | 5140800 | 1.6403e+004 | 13.7731 | 0.137417 |
| 5 | 981x1470x3 | 326210 | 1.6161e+004 | 13.9219 | 0.118353 |
| 6 | 165x247x3 | 122265 | 1.1708e+004 | 17.1454 | 0.007986 |
| 7 | 360x480x3 | 518400 | 2.3129e+004 | 10.3367 | 0.018330 |
| 8 | 183x275x3 | 150975 | 2.1813e+004 | 10.9226 | 0.006040 |
| 9 | 183x275x3 | 150975 | 1.6117e+004 | 13.9491 | 0.006559 |
| 10 | 201x251x3 | 151353 | 2.4935e+004 | 9.5848 | 0.007003 |
| 11 | 600x1050x3 | 18000 | 2.3488e+004 | 10.1830 | 0.060182 |
| 12 | 1144x1783x3 | 6119256 | 6.9825e+003 | 22.3136 | 0.166045 |

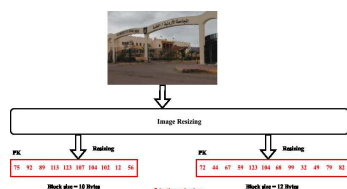


Fig. 5: Selecting the red color of image_key to generate PK

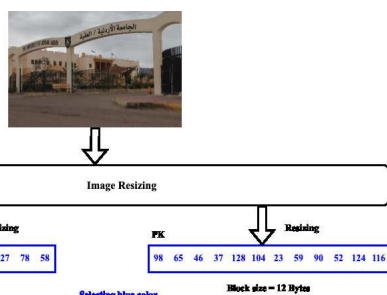


Fig. 6: Selecting the blue color of image_key to generate PK

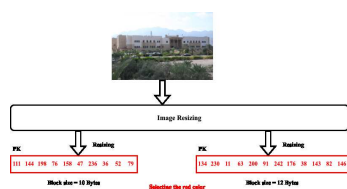


Fig. 7: Selecting the red color of image_key to generate PK

Table 4: Encrypting-decrypting image 2, image 4 is an image_key

| BLS (byte) (number of keys) | MSE | PSNR | Encryption time (second) |
|-----------------------------|--------------------|----------------|--------------------------|
| 16 | 1.5070e+004 | 14.6209 | 0.037891 |
| 32 | 1.5373e+004 | 14.4216 | 0.030081 |
| 40 | 1.5809e+004 | 14.1420 | 0.027030 |
| 100 | 1.5393e+004 | 14.4083 | 0.024358 |
| 500 | 1.5369e+004 | 14.4241 | 0.025259 |
| 800 | 1.5453e+004 | 14.3699 | 0.022982 |
| 1000 | 1.5265e+004 | 14.4923 | 0.024595 |

Table 5: Encrypting-decrypting image 12, image 4 is an image_key

| BLS (byte) (number of keys) | MSE | PSNR | Encryption time (second) |
|-----------------------------|--------------------|----------------|--------------------------|
| 16 | 1.2623e+004 | 16.3927 | 0.857091 |
| 32 | 1.1618e+004 | 17.2218 | 0.502785 |
| 40 | 1.2043e+004 | 16.8631 | 0.382784 |
| 100 | 1.1763e+004 | 17.0983 | 0.175020 |
| 500 | 1.1727e+004 | 17.1285 | 0.074264 |
| 800 | 1.1741e+004 | 17.1172 | 0.064809 |
| 1000 | 1.1688e+004 | 17.1622 | 0.065107 |

Table 6: Encrypting-decrypting image 3, image 4 is an image_key

| BLS (byte) (number of keys) | MSE | PSNR | Encryption time (second) |
|-----------------------------|--------------------|----------------|--------------------------|
| 16 | 1.3756e+004 | 15.5326 | 0.094956 |
| 32 | 1.3186e+004 | 15.9563 | 0.059891 |
| 40 | 1.3521e+004 | 15.7051 | 0.058479 |
| 100 | 1.3362e+004 | 15.8237 | 0.035774 |
| 500 | 1.3358e+004 | 15.8265 | 0.028064 |
| 800 | 1.3358e+004 | 15.8268 | 0.026850 |
| 1000 | 1.3315e+004 | 15.8589 | 0.029128 |

In addition, we can observe that the proposed method can be easily implemented in parallel using a multithreading system with various threads. There is a significant speedup using two or more threads. Also, the proposed

method satisfies the image quality requirement by providing excellent value for MSE and PSNR in both the encryption and decryption phases. The proposed method

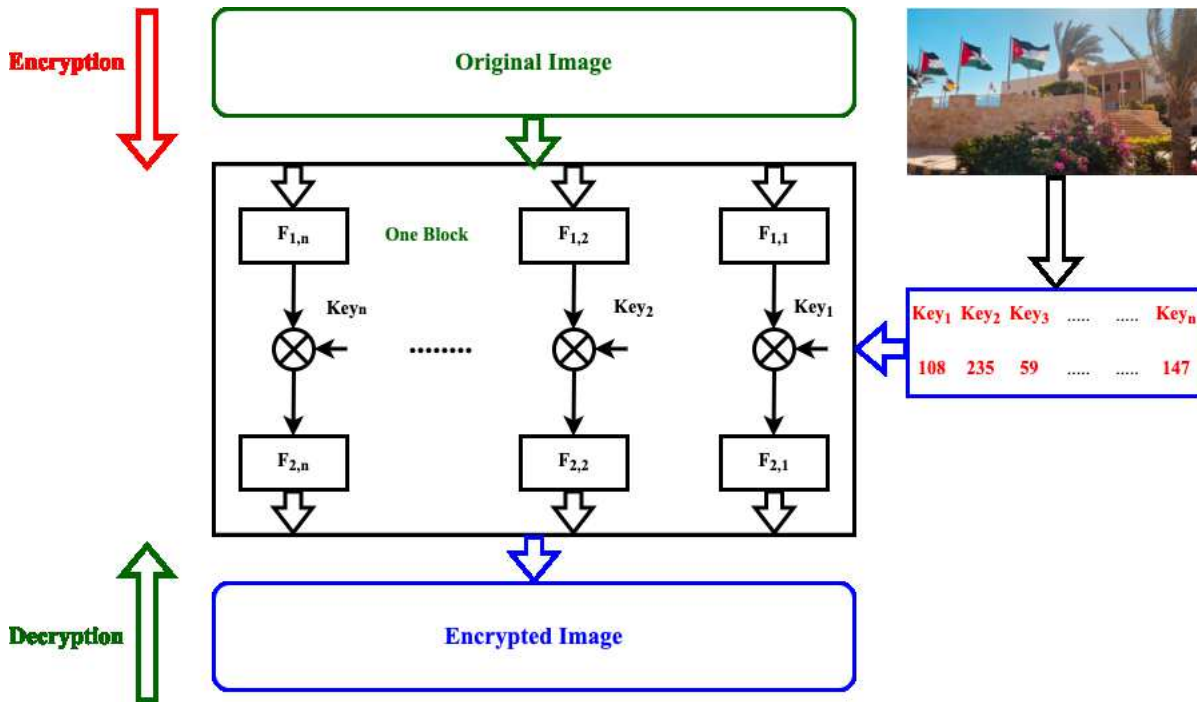


Fig. 8: Selecting the blue color of another image_key to generate PK



Fig. 9: Selecting the blue color of another image_key to generate PK

```

matlabpool('open', 4)
par for i=1:r:n1*n2*n3-r+1
.
.
end
matlabpool close
    
```

In this case, the computational operation is divided into 4 parts (threads), which are dealt with separately from respective core processors. These computing tasks take place dynamically when the beginning of the next cycle step assigns the proper computer operations to free the processor core, leading to optimal utilization.

Fig. 10: Multithreading Matlab pool

The proposed method results were compared with the methods proposed in [30,31,32,33,34], and the proposed method show a significant speedup as show in Table 9:

The proposed method adds the following improvements to the standard techniques of data cryptography:

- 1.It maintains a higher security level. The private key length is variable
- 2.The number of generated keys will equal the block size; the private key contents depend on the selected image_key and the block size.
- 3.One round is required, and expanding the number of rounds is straightforward.
- 4.Using the proposed approach would speed up the cryptography process.

can be used for various data types, including color images and text files.

Table 7: Methods comparisons

| Image size(byte) | Encryption time(second) | | | | |
|---|-------------------------|----------------|----------------|----------------|--|
| | DES | 3DES | AES | BF | MPK (BLS=16 byte=128 bits, rounds=10) |
| 150849 | 0.1093 | 0.1249 | 0.1008 | 0.0603 | 0.009153 |
| 77976 | 0.0572 | 0.0647 | 0.0522 | 0.0313 | 0.004588 |
| 518400 | 0.3761 | 0.4292 | 0.3456 | 0.2067 | 0.016747 |
| 5140800 | 3.6998 | 4.2546 | 3.4248 | 2.0452 | 0.137417 |
| 4326210 | 3.1138 | 3.5808 | 2.8822 | 1.7213 | 0.118353 |
| 122265 | 0.0884 | 0.1015 | 0.0818 | 0.0490 | 0.007986 |
| 518400 | 0.3739 | 0.4294 | 0.3457 | 0.2064 | 0.018330 |
| 150975 | 0.1085 | 0.1253 | 0.1012 | 0.0607 | 0.006040 |
| 150975 | 0.1089 | 0.1252 | 0.1012 | 0.0607 | 0.006559 |
| 151353 | 0.1092 | 0.1257 | 0.1013 | 0.0604 | 0.007003 |
| 1890000 | 1.3607 | 1.5646 | 1.2595 | 0.7521 | 0.060182 |
| 6119256 | 4.4038 | 5.0648 | 4.0767 | 2.4345 | 0.166045 |
| Average | 1.1591 | 1.3326 | 1.0728 | 0.6407 | 0.0465 |
| Throughput (K byte) | 1356.3 | 1179.7 | 1465.4 | 2453.7 | 33808 |
| Speedup of the Proposed method | 24.9266 | 28.6581 | 23.0708 | 13.7784 | 1.0000 |

Table 8: Proposed method implementation on multithreading environment

| Image number | 1 thread time | 2 threads time | 4 threads time | 8 threads time |
|----------------|---------------|----------------|----------------|----------------|
| 1 | 0.044650 | 0.0263 | 0.0123 | 0.0068 |
| 2 | 0.037891 | 0.0234 | 0.0105 | 0.0058 |
| 3 | 0.094956 | 0.0519 | 0.0259 | 0.0142 |
| 4 | 0.767682 | 0.3998 | 0.2069 | 0.1142 |
| 5 | 0.614167 | 0.3232 | 0.1660 | 0.0915 |
| 6 | 0.043193 | 0.0259 | 0.0119 | 0.0066 |
| 7 | 0.098163 | 0.0536 | 0.0267 | 0.0147 |
| 8 | 0.044476 | 0.0254 | 0.0122 | 0.0067 |
| 9 | 0.042990 | 0.0246 | 0.0118 | 0.0064 |
| 10 | 0.042769 | 0.0244 | 0.0117 | 0.0064 |
| 11 | 0.299206 | 0.1662 | 0.0818 | 0.0447 |
| 12 | 0.879196 | 0.4440 | 0.2345 | 0.1303 |
| Average | 0.2508 | 0.1324 | 0.0677 | 0.0374 |
| Speedup | 1.0000 | 1.8943 | 3.7046 | 6.7059 |

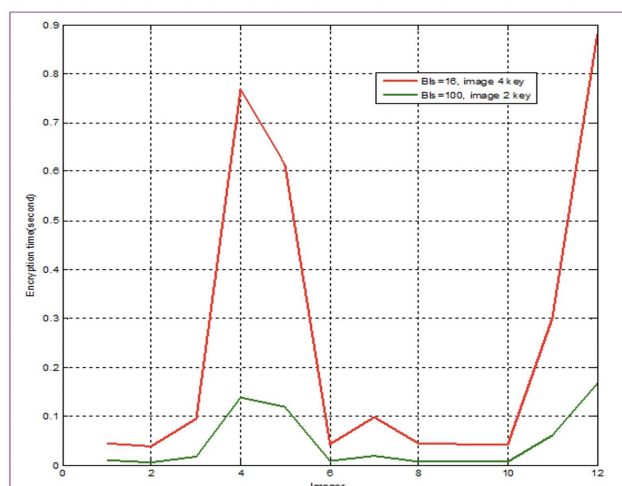


Fig. 11: The relationship between Encryption time and image size

6 Conclusions

A method of color image cryptography was introduced, programmed, and executed. This method can encrypt-decrypt any data, including color images. The proposed method provides a high-security level; this can be achieved depending on the selected image_key, the

color matrix, and the data block size, which form the length and contents of the secret private key. It was shown that the private key is not fixed and changed dynamically when the image key or the selected color and data block

Table 9: Throughput comparisons

| Method | Throughput (K bytes per second) | Speed up of the proposed method |
|--|------------------------------------|------------------------------------|
| Ref. [21] | 888.8867 | 40.4877 |
| Ref. [22] | 638.4082 | 56.3730 |
| Ref. [23] | 911.0352 | 39.5034 |
| Ref. [24] | 361.4102 | 99.5794 |
| Ref. [25] | 384.9609 | 93.4874 |
| Proposed* | 35989 | 1.0000 |
| For image 12 throughput= 6119256/(0.166045*1024)= 35989 K bytes per second | | |

size changed. The proposed method enhances the data cryptography process, providing a good speedup compared with other standard methods; this speedup can be increased by implementing the proposed method using a multithreading environment. The obtained MSE and PSNR values were acceptable for both the encryption and decryption phases. The proposed method can be in the future implemented using hardware; here, a particular processor can be efficiently designed to handle the operations of the proposed method.

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