



Characterizing affine \mathcal{C} -semigroups

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Abstract

Let $\mathcal{C} \subset \mathbb{N}^p$ be a finitely generated integer cone and $S \subset \mathcal{C}$ be an affine semigroup such that the real cones generated by \mathcal{C} and by S are equal. The semigroup S is called \mathcal{C} -semigroup if $\mathcal{C} \setminus S$ is a finite set. In this paper, we characterize the \mathcal{C} -semigroups from their minimal generating sets, and we give an algorithm to check if S is a \mathcal{C} -semigroup and to compute its set of gaps. We also study the embedding dimension of \mathcal{C} -semigroups obtaining a lower bound for it, and introduce some families of \mathcal{C} -semigroups whose embedding dimension reaches our bound. In the last section, we present a method to obtain a decomposition of a \mathcal{C} -semigroup into irreducible \mathcal{C} -semigroups.

Keywords Affine semigroup · \mathcal{C} -semigroup · Embedding dimension · Gap of a semigroup · Generalized numerical semigroup · Irreducible semigroup

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1 Introduction

An affine semigroup $S \subset \mathbb{N}^p$ is called \mathcal{C}_S -semigroup if $\mathcal{C}_S \setminus S$ is a finite set where $\mathcal{C}_S \subset \mathbb{N}^p$ is the minimal integer cone containing it. These semigroups are a natural generalization of numerical semigroups, and several of their invariants can be generalized. For a given numerical semigroup G , it is well-known that $\mathbb{N} \setminus G$ is finite; in fact, $G \subset \mathbb{N}$ is a numerical semigroup if it is a submonoid of \mathbb{N} and $\mathbb{N} \setminus G$ is finite (for topics related with numerical semigroups see [13] and the references therein). In general, it does not happen for affine semigroups.

\mathcal{C} -semigroups are introduced in [8], where the authors study several properties about them (for example, an extended Wilf's conjecture for \mathcal{C} -semigroups is given). These semigroups appear in different contexts: when the integer points in an infinite family of some homothetic convex bodies in \mathbb{R}_{\geq}^p are considered (see, for instance, [9], [10] and the references therein), or when the non-negative integer solutions of some modular Diophantine inequality are studied (see [5]), et cetera. In case the cone \mathcal{C} is \mathbb{N}^p , \mathbb{N}^p -semigroups are called generalized numerical semigroups and they were introduced in [6]. Recently, in [11] it is proved that the minimal free resolution of the associated algebra to any \mathcal{C} -semigroup has maximal projective dimension possible.

In this context, \mathbb{N}^p -semigroups are characterized in [3], but the general problem was opened, *given any affine semigroup S , how to detect if S is or not a \mathcal{C}_S -semigroup?* The primary goal of this work is to determine the conditions that any affine semigroup given by its minimal set of generators has to verify to be a \mathcal{C}_S -semigroup. We solve this problem in Theorem 9, and in Algorithm 1 we provide a computational way to check it.

Another open problem is to compute the set of gaps of any \mathcal{C} -semigroup defined by its minimal generating set. We solve this problem by means of setting a finite subset of \mathcal{C} containing all the gaps of a given \mathcal{C} -semigroup. Algorithm 2 computes the set of gaps of the given \mathcal{C} -semigroup.

In this paper, we also go in-depth to study the embedding dimension of \mathcal{C} -semigroups. In [8, Theorem 11], a lower bound of the embedding dimension of \mathbb{N}^p -semigroups is provided, and some families of \mathbb{N}^p -semigroups reaching this bound are given. Besides, in [8, Conjecture 12], it is proposed a conjecture about a lower bound for the embedding dimension of any \mathcal{C} -semigroup. In Sect. 5, we introduce a lower bound of the embedding dimension of any \mathcal{C} -semigroup, and some families of \mathcal{C} -semigroups whose embedding dimension is equal to this new bound.

An important problem in Semigroup Theory is to determine some decomposition of a semigroup into irreducible semigroups (for example, see [13, Chapter 3] for numerical semigroups, or its generalization for \mathbb{N}^p -semigroups in [2]). We propose an algorithm to compute a decomposition of any \mathcal{C} -semigroup into irreducible \mathcal{C} -semigroups.

The results of this work are illustrated with several examples. To this aim, we have used third-party software, such as Normaliz [4], and the libraries `CharacterizingAffineCSemigroup` and `Irreducible` [7] developed by the authors in Python [12].

The content of this work is organized as follows. Section 2 introduces the initial definitions and notations used throughout the paper, mainly related to finitely generated cones. In Sect. 3, a characterization of \mathcal{C} -semigroups is provided, and an algorithm

to check if an affine semigroup is a \mathcal{C} -semigroup. Section 4 is devoted to give an algorithm to compute the set of gaps of a \mathcal{C} -semigroup. Section 5 makes a study of the minimal generating sets of \mathcal{C} -semigroups formulating explicitly a lower bound for their embedding dimensions. Finally, in Sect. 6 an algorithm for computing a decomposition of a \mathcal{C} -semigroup into irreducible \mathcal{C} -semigroups is presented.

2 Preliminaries

The sets of real numbers, rational numbers, integer numbers and the non-negative integer numbers are denoted by \mathbb{R} , \mathbb{Q} , \mathbb{Z} and \mathbb{N} , respectively. Given a subset A of \mathbb{R} , A_{\geq} is the set of elements in A greater than or equal to zero. For any $n \in \mathbb{N}$, $[n]$ denotes the set $\{1, \dots, n\}$. Given an element x in \mathbb{R}^n , $\|x\|_1$ denotes the sum of the absolute value of its entries, that is, its 1-norm. In this paper we assume the set $\{\mathbf{e}_1, \dots, \mathbf{e}_p\}$ is the canonical basis of \mathbb{R}^p .

For a non empty subset B of \mathbb{R}_{\geq}^p , we define the cone generated by B :

$$L(B) := \left\{ \sum_{i=1}^n \lambda_i \mathbf{b}_i \mid n \in \mathbb{N}, \{\mathbf{b}_1, \dots, \mathbf{b}_n\} \subset B, \text{ and } \lambda_i \in \mathbb{R}_{\geq}, \forall i \in [n] \right\}.$$

Given a real cone $\mathcal{C} \subset \mathbb{R}_{\geq}^p$, it is well-known that $\mathcal{C} \cap \mathbb{N}^p$ is finitely generated if and only if there exists a rational point in each extremal ray of \mathcal{C} . Moreover, any subsemigroup of \mathcal{C} is finitely generated if and only if there exists an element in the semigroup in each extremal ray of \mathcal{C} . A good monograph about rational cones and affine monoids is [1]. From now on, we assume that the integer cones considered in this work are finitely generated.

Definition 1 Given an integer cone $\mathcal{C} \subset \mathbb{N}^p$, an affine semigroup $S \subset \mathcal{C}$ is said to be a \mathcal{C} -semigroup if $\mathcal{C} \setminus S$ is a finite set. If the cone $\mathcal{C} = \mathbb{N}^p$, a \mathcal{C} -semigroup is called \mathbb{N}^p -semigroup.

Fix a finitely generated semigroup $S \subset \mathbb{N}^p$, we denote by \mathcal{C}_S the integer cone $L(S) \cap \mathbb{N}^p$. Note that, if S is a \mathcal{C} -semigroup, the cone \mathcal{C} is \mathcal{C}_S . Obviously, a unique cone corresponds to infinite different semigroups.

The cone $L(S)$ is a polyhedron and we denote by $\{h_1(x) = 0, \dots, h_t(x) = 0\}$ the set of its supported hyperplanes. We suppose $L(S) = \{x \in \mathbb{R}_{\geq}^d \mid h_1(x) \geq 0, \dots, h_t(x) \geq 0\}$. Unless otherwise stated, the considered coefficients of each $h_i(x)$ are integers and relatively primes.

Assume $L(S)$ has q extremal rays denoted by τ_1, \dots, τ_q . Then, each τ_i is determined by the set of linear equations $H_i := \{h_{j_1^{(i)}}(x) = 0, \dots, h_{j_{p-1}^{(i)}}(x) = 0\}$ where $J_i := \{j_1^{(i)} < \dots < j_{p-1}^{(i)}\} \subset [t]$ is the index set of the supported hyperplanes containing τ_i . So, for each $i \in [q]$, there exists the minimal non-negative integer vector \mathbf{a}_i such that $\tau_i = \{\lambda \mathbf{a}_i \mid \lambda \in \mathbb{R}_{\geq}\}$. The set $\{\mathbf{a}_1, \dots, \mathbf{a}_q\}$ is a generating set of $L(S)$.

Note that a necessary condition for S to be a \mathcal{C}_S -semigroup is the set $\tau_i \cap (\mathcal{C}_S \setminus S)$ is finite for all $i \in [q]$.

From each extremal ray τ_i of $L(S)$, we define $v_i(\alpha)$ as the parallel line to τ_i given by the solutions of the linear equations $\bigcup_{j \in J_i} \{h_j(x) = \alpha_j\}$ where $\alpha = (\alpha_{j_1}, \dots, \alpha_{j_{p-1}}) \in \mathbb{Z}^{p-1}$. For every integer point $P \in \mathbb{Z}^p$ and $i \in [q]$, there exists $\alpha \in \mathbb{Z}^{p-1}$ such that P belongs to $v_i(\alpha)$; if $P \in \mathcal{C}_S$, $\alpha \in \mathbb{N}^{p-1}$. We denote by $\Upsilon_i(P)$ the element $(h_{j_1^{(i)}}(P), \dots, h_{j_{p-1}^{(i)}}(P)) \in \mathbb{N}^{p-1}$ with $J_i = \{j_1^{(i)} < \dots < j_{p-1}^{(i)}\}$, $P \in \mathcal{C}_S$ and $i \in [q]$. Note that for any $P \in \mathcal{C}_S$, $P \in v_i(\alpha)$ if and only if $\alpha = \Upsilon_i(P)$.

Since all the semigroups appearing in this work are finitely generated, from now on, we omit the term *affine* when affine semigroups are considered.

3 An algorithm to detect if a semigroup is a \mathcal{C} -semigroup

In this section, we study the conditions that a semigroup has to satisfy to be a \mathcal{C} -semigroup. This characterization depends on the minimal set of generators of the given semigroup.

Let $S \subset \mathbb{N}^p$ be the affine semigroup minimally generated by $\Lambda_S = \{\mathbf{s}_1, \dots, \mathbf{s}_q, \mathbf{s}_{q+1}, \dots, \mathbf{s}_n\}$ and τ_1, \dots, τ_q be the extremal rays of $L(S)$. Assume that for every $i \in [q]$, $\tau_i \cap (\mathcal{C}_S \setminus S)$ is finite and \mathbf{s}_i is the minimum (respect to the natural order) element in Λ_S belonging to τ_i . We denote by \mathbf{f}_i the maximal element in $\tau_i \cap (\mathcal{C}_S \setminus S)$ with respect to the natural order in \mathbb{N}^p . Recall that \mathbf{a}_i is the minimal non-negative integer vector defining τ_i , and let $\mathbf{c}_i \in S$ be the element $\mathbf{f}_i + \mathbf{a}_i$. In case $\tau_i \cap (\mathcal{C}_S \setminus S) = \emptyset$, we fix $\mathbf{f}_i = -\mathbf{a}_i$. The elements \mathbf{f}_i and \mathbf{c}_i are a generalization on the semigroup $\tau_i \cap S$ of the concepts Frobenius number and conductor of a numerical semigroup; for numerical semigroups, the Frobenius number is the maximal natural number that is not in the semigroup, and the conductor is Frobenius number plus one (see [13, Chapter 1]). Hence, we call Frobenius element and conductor of the semigroup $\tau_i \cap S$ the elements \mathbf{f}_i and \mathbf{c}_i , respectively. One easy but important property of S is for every $P \in S$, $P + \mathbf{c}_i + \lambda \mathbf{a}_i \in S$ for any $i \in [q]$ and $\lambda \in \mathbb{N}$.

Note that $\tau_i \cap \mathbb{N}^p$ is equal to $\{\lambda \mathbf{a}_i \mid \lambda \in \mathbb{N}\}$. So, there exists $S_i \subset \mathbb{N}$ such that $\tau_i \cap S = \{\lambda \mathbf{a}_i \mid \lambda \in S_i\}$. If we assume that $\tau_i \cap (\mathcal{C}_S \setminus S)$ is finite, it is easy to prove that S_i is a numerical semigroup.

Lemma 2 *The τ_i -semigroup $\tau_i \cap S$ is isomorphic to the semigroup $S_i = \{\lambda \in \mathbb{N} \mid \lambda \mathbf{a}_i \in S\}$. Moreover $\tau_i \cap (\mathcal{C}_S \setminus S)$ is finite if and only if S_i is a numerical semigroup.*

Proof Consider the isomorphism $\varphi : \tau_i \cap S \rightarrow S_i$ with $\varphi(\mathbf{w}) := \lambda$ such that $\mathbf{w} = \lambda \mathbf{a}_i$. The second statement holds since $\tau_i \cap (\mathcal{C}_S \setminus S) = \{\lambda \mathbf{a}_i \in \mathbb{N} \mid \lambda \mathbf{a}_i \notin S, \lambda \in \mathbb{N}\}$. \square

Corollary 3 *Given the semigroup $\tau_i \cap S$, \mathbf{f}_i is equal to $f \mathbf{a}_i$ and $\mathbf{c}_i = c \mathbf{a}_i$ where f and c are the Frobenius number and the conductor of the numerical semigroup S_i , respectively.*

To test whether $\tau_i \cap (\mathcal{C}_S \setminus S)$ is finite, the following result can be used.

Lemma 4 *Let $S \subset \mathbb{N}^p$ be a semigroup and τ be an extremal ray of $L(S)$ satisfying $\tau \cap \mathbb{N}^p = \{\lambda \mathbf{a} \mid \lambda \in \mathbb{N}\}$ with $\mathbf{a} \in \mathbb{N}^p$. Then, $\tau \cap (\mathcal{C}_S \setminus S)$ is finite if and only if $\gcd(\{\lambda \mid \lambda \mathbf{a} \in \tau \cap \Lambda_S\}) = 1$.*

Proof Assume that $\tau \cap (\mathcal{C}_S \setminus S)$ is finite and suppose that $\gcd(\{\lambda \mid \lambda \mathbf{a} \in \tau \cap \Lambda_S\}) = n \neq 1$. Hence, every element $\lambda \mathbf{a}$ with $\gcd(n, \lambda) = 1$ does not belong to S , and then $\tau \cap (\mathcal{C}_S \setminus S)$ is not finite.

Conversely if $\gcd(\{\lambda \mid \lambda \mathbf{a} \in \tau \cap \Lambda_S\}) = 1$, then the semigroup $S' = \{\lambda \in \mathbb{N} \mid \lambda \mathbf{a} \in S\}$ is a numerical semigroup. From the proof of Lemma 2, S' is isomorphic to $\tau \cap S$. Therefore, $\tau \cap (\mathcal{C}_S \setminus S)$ is finite. \square

To introduce the announced characterization, we need to define some subsets of $L(S)$ and prove some of their properties. Associated to the integer cone \mathcal{C}_S , consider the sets $\mathcal{A} := \{\sum_{i \in [q]} \lambda_i \mathbf{a}_i \mid 0 \leq \lambda_i \leq 1\} \cap \mathbb{N}^p$ and $\mathcal{D} := \{\sum_{i \in [q]} \lambda_i \mathbf{s}_i \mid 0 \leq \lambda_i \leq 1\} \cap \mathbb{N}^p$.

Lemma 5 *Given $P \in \mathcal{C}_S$, there exist $Q \in \mathcal{A}$ and $\beta \in \mathbb{N}^q$ such that $P = Q + \sum_{i \in [q]} \beta_i \mathbf{a}_i$. Moreover, $\Upsilon_j(P) = \Upsilon_j(Q) + \sum_{i \in [q]} \beta_i \Upsilon_j(\mathbf{a}_i)$ for every $j \in [q]$.*

Proof Since $P \in \mathcal{C}_S$, $P = \sum_{i \in [q]} \mu_i \mathbf{a}_i$ with $\mu_i \in \mathbb{Q}_{\geq}$. For each μ_i there exists $\lambda_i \in [0, 1)$ satisfying $\mu_i = \lfloor \mu_i \rfloor + \lambda_i$. Hence, $P = Q + \sum_{i \in [q]} \lfloor \mu_i \rfloor \mathbf{a}_i$ where $Q = \sum_{i \in [q]} \lambda_i \mathbf{a}_i = P - \sum_{i \in [q]} \lfloor \mu_i \rfloor \mathbf{a}_i \in \mathcal{A}$. Trivially, $\Upsilon_j(P)$ is equal to $\Upsilon_j(Q) + \sum_{i \in [q]} \beta_i \Upsilon_j(\mathbf{a}_i)$ for every $j \in [q]$. \square

For every $i \in [q]$, consider $T_i \subset \mathbb{N}^{p-1}$ the semigroup generated by the finite set $\{\Upsilon_i(Q) \mid Q \in \mathcal{A}\}$ and let Γ_i be its minimal generating set. Note that the sets \mathcal{A}, T_i and Γ_i only depend on the cone \mathcal{C}_S , and $0 \in T_i$, since $\mathbf{a}_i \in \mathcal{A}$. The relationships between the elements in \mathcal{C}_S and S , and the elements belonging to T_i and Γ_i are explicitly determined in the following results for each $i \in [q]$.

Lemma 6 *Let P be an element in \mathcal{C}_S such that $P \in v_i(\alpha)$ for some $\alpha \in \mathbb{N}^{p-1}$, then $\alpha \in T_i$.*

Proof By definition, $P \in v_i(\alpha)$ means that $\alpha = \Upsilon_i(P)$. Using Lemma 5, $P = Q + \sum_{j \in [q]} \beta_j \mathbf{a}_j$ with $Q, \mathbf{a}_1, \dots, \mathbf{a}_q \in \mathcal{A}$ and $\beta_1, \dots, \beta_q \in \mathbb{N}$. Therefore, $\Upsilon_i(P) = \Upsilon_i(Q) + \sum_{j \in [q]} \beta_j \Upsilon_i(\mathbf{a}_j) \in T_i$. \square

Corollary 7 *For every $\alpha \in T_i$, $\mathcal{C}_S \cap v_i(\alpha) \neq \emptyset$ if and only if $\mathcal{C}_S \cap v_i(\beta) \neq \emptyset$ for all $\beta \in \Gamma_i$.*

Proof Since $\Gamma_i \subset T_i$, if for all $\alpha \in T_i$, $\mathcal{C}_S \cap v_i(\alpha) \neq \emptyset$ then $\mathcal{C}_S \cap v_i(\beta) \neq \emptyset$ for all $\beta \in \Gamma_i$.

Assume that $\mathcal{C}_S \cap v_i(\beta) \neq \emptyset$ for all $\beta \in \Gamma_i$ and let α be an element in T_i . Then, there exist $\beta_1, \dots, \beta_k \in \Gamma_i, \mu_1, \dots, \mu_k \in \mathbb{N}$ and $Q_1, \dots, Q_k \in \mathcal{A}$ such that $\alpha = \sum_{j \in [k]} \mu_j \beta_j$ and $\Upsilon_i(Q_j) = \beta_j$ for $j \in [k]$. Note that $P = \sum_{j \in [k]} \mu_j Q_j \in \mathcal{C}_S$ belongs to $v_i(\alpha)$. \square

Corollary 8 *For every $\alpha \in T_i$, $S \cap v_i(\alpha) \neq \emptyset$ if and only if $S \cap v_i(\beta) \neq \emptyset$ for all $\beta \in \Gamma_i$.*

Proof Since $S \cap v_i(\beta) \neq \emptyset$ for all β , then there exists $Q_1, \dots, Q_k \in \mathcal{A}$ such that $\Upsilon_i(Q_j) = \beta_j$ for $j \in [k]$. Thus, the proof of this corollary is analogous to the proof of Corollary 7. \square

Note that if $P \in S \cap v_i(\alpha)$ for some $\alpha \in \mathbb{N}^{p-1}$ and $i \in [q]$, then $P + \mathbf{c}_i + \lambda \mathbf{a}_i \in S$ and $\Upsilon_i(P + \mathbf{c}_i + \lambda \mathbf{a}_i) = \alpha$ for all $\lambda \in \mathbb{N}$.

Now, we introduce a characterization of \mathcal{C} -semigroups. This characterization depends on the minimal generating set of the given semigroup. Besides, from its proof, we provide an algorithm for checking if a semigroup is a \mathcal{C} -semigroup (Algorithm 1). Note that most of the parts of Algorithm 1 can be parallelized at least in q stand-alone processes.

Theorem 9 *A semigroup S minimally generated by $\Lambda_S = \{\mathbf{s}_1, \dots, \mathbf{s}_n\}$ is a \mathcal{C}_S -semigroup if and only if:*

1. $\tau_i \cap (\mathcal{C}_S \setminus S)$ is finite for all $i \in [q]$.
2. $\Lambda_S \cap v_i(\alpha) \neq \emptyset$ for all $\alpha \in \Gamma_i$ and $i \in [q]$.

Proof Let S be a \mathcal{C}_S -semigroup. Trivially, $\tau_i \cap (\mathcal{C}_S \setminus S)$ is finite for all $i \in [q]$. Now let $i \in [q]$ and $\alpha \in \Gamma_i$, we prove that $\Lambda_S \cap v_i(\alpha) \neq \emptyset$. Since $\alpha \in \Gamma_i$, there exists $Q \in \mathcal{A}$ such that $\alpha = \Upsilon_i(Q)$. Besides, $Q + \lambda \mathbf{a}_i \in \mathcal{C}_S$ and $\Upsilon_i(Q + \lambda \mathbf{a}_i) = \alpha$ for all $\lambda \in \mathbb{N}$. For some $\lambda \in \mathbb{N}$, $Q + \lambda \mathbf{a}_i$ has to be in S (S is \mathcal{C}_S -semigroup), that is to say, $Q + \lambda \mathbf{a}_i = \sum_{j \in [n]} \mu_j \mathbf{s}_j$ with $\mu_1, \dots, \mu_n \in \mathbb{N}$. Therefore, $\alpha = \Upsilon_i(Q + \lambda \mathbf{a}_i) = \sum_{j \in [n]} \mu_j \Upsilon_i(\mathbf{s}_j)$. By Lemma 5, for all $j \in [n]$, $\mathbf{s}_j = Q_j + \sum_{k \in [q]} \beta_{jk} \mathbf{a}_k$ for some $Q_j \in \mathcal{A}$ and $\beta_{j1}, \dots, \beta_{jq} \in \mathbb{N}$. So, $\alpha = \sum_{j \in [n]} \mu_j \Upsilon_i(Q_j + \sum_{k \in [q]} \beta_{jk} \mathbf{a}_k) = \sum_{j \in [n]} \mu_j \Upsilon_i(Q_j) + \sum_{j \in [n]} \sum_{k \in [q]} \mu_j \beta_{jk} \Upsilon_i(\mathbf{a}_k)$. Since α is a minimal generator of T_i , $\sum_{j \in [n]} \mu_j + \sum_{j \in [n]} \sum_{k \in [q] \setminus \{i\}} \mu_j \beta_{jk} = 1$. So $\beta_{jk} = 0$ for all $j \in [n]$ and for all $k \in [q] \setminus \{i\}$, and there exists $l \in [n]$ such that $\mu_l = 1$ and $\mu_j = 0$ for all $j \in [n] \setminus \{l\}$. Hence, there exists $\mathbf{s} \in \Lambda_S$ such that $\Upsilon_i(\mathbf{s}) = \alpha$ and then $\Lambda_S \cap v_i(\alpha) \neq \emptyset$.

Conversely, we assume that $\forall i \in [q]$ and $\forall \alpha \in \Gamma_i$, $\tau_i \cap (\mathcal{C}_S \setminus S)$ is finite and $\Lambda_S \cap v_i(\alpha) \neq \emptyset$ (recall that $\mathbf{c}_i = \mathbf{f}_i + \mathbf{a}_i$). Let Q be an element in \mathcal{D} . By Lemmas 5 and 6, $Q \in \alpha_i(\Upsilon_i(Q))$ and $\Upsilon_i(Q) \in T_i$. If $Q \in \tau_i$ for some $i \in [q]$, then $v_i(\Upsilon_i(Q)) = \tau_i$ and, by the first condition, $S \cap v_i(\Upsilon_i(Q)) \neq \emptyset$. If Q is not in any ray, by the second condition and Corollary 8, $S \cap v_i(\Upsilon_i(Q)) \neq \emptyset$. Therefore, for every $Q \in \mathcal{D}$, the line $v_i(\Upsilon_i(Q))$ includes a unique non zero minimum (respect 1-norm) point belonging to S . Denote by $\{\mathbf{m}_{i1}, \dots, \mathbf{m}_{id_i}\}$ the set obtained from the union of above points for the different elements in \mathcal{D} (some of these elements belong to Λ_S). Note that $\mathbf{m}_{ij} + \mathbf{c}_i + \lambda \mathbf{a}_i \in S$ for all $j \in [d_i]$ and $\lambda \in \mathbb{N}$. Consider $n_i := \max\{\|\mathbf{m}_{i1} + \mathbf{c}_i\|_1, \dots, \|\mathbf{m}_{id_i} + \mathbf{c}_i\|_1\}$, and \mathbf{x}_i the minimum element (respect to the 1-norm) in $\tau_i \cap S$ such that $\|\mathbf{x}_i\|_1$ is greater than or equal to n_i . The set $\mathcal{D}_i := \mathcal{D} + \mathbf{x}_i$ satisfies that $\mathcal{D}_i \cap S = \mathcal{D}_i \cap \mathcal{C}_S = \mathcal{D}_i$. Consider $\mathbf{a} \in \mathbf{x}_i + \mathcal{C}_S$, proceeding as in the proof of Lemma 5, $\mathbf{a} = \mathbf{x}_i + P + \sum_{j \in [q]} \beta_j \mathbf{s}_j$ for some $P \in \mathcal{D}$ and $\beta_1, \dots, \beta_q \in \mathbb{N}$, and hence, $\mathbf{x}_i + \mathcal{C}_S \subset S$. We define the bounded set $\mathcal{X} := \{\sum_{i \in [q]} \lambda_i \mathbf{x}_i \mid 0 \leq \lambda_i \leq 1\}$. Since $\mathbf{x}_i + \mathcal{C}_S \subset S$ for every $i \in [q]$ and $L(\mathcal{C}_S) = \{\sum_{i \in [q]} \lambda_i \mathbf{x}_i \mid \lambda_i \in \mathbb{R}_{\geq}\}$, $\mathcal{C}_S \setminus S \subset \mathcal{X}$. Therefore, S is a \mathcal{C}_S -semigroup. \square

Example 10 illustrates Theorem 9 and Algorithm 1.

Example 10 Let $S \subset \mathbb{N}^3$ be the semigroup minimally generated by

$$\Lambda_S = \{(2, 0, 0), (4, 2, 4), (0, 1, 0), (3, 0, 0), (6, 3, 6), (3, 1, 1), (4, 1, 1), (3, 1, 2), (1, 1, 0), (3, 2, 3), (1, 2, 1)\}.$$

Algorithm 1: Test if a semigroup S is a \mathcal{C}_S -semigroup.

Input: The minimal generating set Λ_S of a semigroup $S \subset \mathbb{N}^P$.
Output: Check if S is a \mathcal{C}_S -semigroup.
begin
 $q \leftarrow$ number of extremal rays of $L(S)$;
 if $\tau_i \cap (\mathcal{C}_S \setminus S)$ is not finite for some $i \in [q]$ **then**
 return S is not a \mathcal{C}_S -semigroup.
 Compute the set $\{\mathbf{a}_1, \dots, \mathbf{a}_q\}$ from $L(S)$;
 $\mathcal{A} \leftarrow \{\sum_{i \in [q]} \lambda_i \mathbf{a}_i \mid 0 \leq \lambda_i \leq 1\} \cap \mathbb{N}^P$;
 forall the $i \in [q]$ **do**
 $\Gamma_i \leftarrow$ the minimal generating set of T_i obtained from the finite set $\Upsilon_i(\mathcal{A})$;
 if $\Lambda_S \cap \nu_i(\alpha) \neq \emptyset$ for all $\alpha \in \Gamma_i$ and $i \in [q]$ **then**
 return S is a \mathcal{C}_S -semigroup.
 return S is not a \mathcal{C}_S -semigroup.

The cone $L(S)$ is $\langle (1, 0, 0), (2, 1, 2), (0, 1, 0) \rangle_{\mathbb{R}_\geq}$ and its supported hyperplanes are $h_1(x, y, z) \equiv 2y - z = 0$, $h_2(x, y, z) \equiv x - z = 0$ and $h_3(x, y, z) \equiv z = 0$. Recall $\mathcal{C}_S = L(S) \cap \mathbb{N}^3$. By \mathbf{a}_1 , \mathbf{a}_2 and \mathbf{a}_3 we denote the vectors $(1, 0, 0)$, $(2, 1, 2)$ and $(0, 1, 0)$ respectively, and τ_1 , τ_2 and τ_3 are the extremal rays with sets of defining equations $\{h_1(x, y, z) = 0, h_3(x, y, z) = 0\}$, $\{h_1(x, y, z) = 0, h_2(x, y, z) = 0\}$ and $\{h_2(x, y, z) = 0, h_3(x, y, z) = 0\}$, respectively. Hence, $S_1 = (\tau_1 \setminus \{(1, 0, 0)\}) \cap \mathbb{N}^3$, $S_2 = \tau_2 \setminus \{(2, 1, 2)\} \cap \mathbb{N}^3$ and $S_3 = \tau_3 \cap \mathbb{N}^3$, and the first condition in Theorem 9 holds.

The set \mathcal{A} is equal to

$$\{(0, 0, 0), (0, 1, 0), (1, 0, 0), (1, 1, 0), (1, 1, 1), (2, 1, 1), (2, 1, 2), (2, 2, 2), (3, 1, 2), (3, 2, 2)\}, \tag{1}$$

and

$$\begin{aligned} \Upsilon_1(\mathcal{A}) &= \{(0, 0), (0, 2), (1, 1), (2, 0), (2, 2)\}, \\ \Upsilon_2(\mathcal{A}) &= \{(0, 0), (0, 1), (1, 0), (1, 1), (2, 0), (2, 1)\}, \\ \Upsilon_3(\mathcal{A}) &= \{(0, 0), (0, 1), (0, 2), (1, 0), (1, 1), (1, 2)\}. \end{aligned}$$

Therefore, $\Gamma_1 = \{(0, 2), (1, 1), (2, 0)\}$ and $\Gamma_2 = \Gamma_3 = \{(0, 1), (1, 0)\}$.

Since $\Upsilon_1(\{(3, 1, 1), (3, 1, 2), (1, 1, 0)\}) = \Gamma_1$, $\Upsilon_2(\{(3, 1, 2), (3, 2, 3)\}) = \Gamma_2$, and $\Upsilon_3(\{(1, 1, 0), (1, 2, 1)\}) = \Gamma_3$, S satisfies the second condition in Theorem 9. Hence, S is a \mathcal{C}_S -semigroup.

By using our implementation of Algorithm 1, we can confirm that S is a \mathcal{C}_S -semigroup,

```
In [1]: IsCsemigroup([[2, 0, 0], [4, 2, 4], [0, 1, 0], [3, 0, 0],
                    [6, 3, 6], [3, 1, 1], [4, 1, 1], [3, 1, 2], [1, 1, 0], [3, 2, 3],
                    [1, 2, 1]])
Out[1]: True
```

To finish this section, it should be pointed out that there exist some special cases of semigroups where Theorem 9 can be simplified: \mathbb{N}^p -semigroups and two-dimensional case.

Note that, if the integer cone \mathcal{C}_S is \mathbb{N}^p , its supported hyperplanes are $\{x_1 = 0, \dots, x_p = 0\}$. Moreover, since its extremal rays are the axes, $\tau_i \equiv \{\lambda \mathbf{e}_i \mid \lambda \in \mathbb{Q}_{\geq}\}$ is determined by the equations $\cup_{j \in [p] \setminus \{i\}} \{x_j = 0\}$, and for any canonical generator \mathbf{e} of \mathbb{N}^{p-1} , there exists P in \mathbb{N}^p such that $\Upsilon_i(P) = \mathbf{e}$. Furthermore, $\cup_{j \in [p] \setminus \{i\}} \{\Upsilon_i(\mathbf{e}_j)\}$ is the canonical basis of \mathbb{N}^{p-1} . Hence, $\Gamma_1 = \dots = \Gamma_p$ is the canonical basis of \mathbb{N}^{p-1} . From previous considerations, the same characterization of \mathbb{N}^p -semigroups in [3, Theorem 2.8] is obtained from Theorem 9.

Corollary 11 *A semigroup S minimally generated by Λ_S is an \mathbb{N}^p -semigroup if and only if:*

1. for all $i \in [p]$, the non null entries of the elements in $\tau_i \cap \Lambda_S$ are coprime, or $\mathbf{s}_i = \mathbf{e}_i$.
2. for all $i, j \in [p]$ with $i \neq j$, $\mathbf{e}_i + \lambda_j \mathbf{e}_j \in \Lambda_S$ for some $\lambda_j \in \mathbb{N}$.

Focus on two dimensional case, note that the extremal rays and the supported hyperplanes of a cone are equal. Since for each extremal ray the coefficients of its defining linear equation are relatively primes, the linear equations $h_1(x, y) = 1$ and $h_2(x, y) = 1$ always have non-negative integer solutions. So, any semigroup $S \subset \mathbb{N}^2$ is a \mathcal{C}_S -semigroup if and only if $\tau_i \cap (\mathcal{C}_S \setminus S)$ is finite for $i = 1, 2$, and both sets $\Lambda_S \cap \{h_1(x, y) = 1\}$ and $\Lambda_S \cap \{h_2(x, y) = 1\}$ are non empty.

4 Set of gaps of \mathcal{C} -semigroups

This section gives an algorithm to compute the set of gaps of a \mathcal{C} -semigroup, i.e. the set $\mathcal{H}(S) = \mathcal{C}_S \setminus S$. This algorithm is obtained from Theorem 9. To introduce such an algorithm, let us start by redefining some objects used to prove that theorem.

Given S a \mathcal{C}_S -semigroup with q extremal rays, for any $i \in [q]$, let \mathbf{c}_i be the conductor of the semigroup $\tau_i \cap S$. By Corollary 8, for any $\alpha \in \Upsilon_i(\mathcal{D})$ the intersection $\nu_i(\alpha) \cap S$ is not empty. Hence, set $\mathbf{m}_\alpha^{(i)}$ the element in $\nu_i(\alpha) \cap S$ with minimal 1-norm and $\alpha \in \Upsilon_i(\mathcal{D}) \setminus \{0\}$. Note that $\mathbf{m}_\alpha^{(i)} + \mathbf{c}_i + \lambda \mathbf{a}_i \in S$ for all $\lambda \in \mathbb{N}$. Let $n_i := \|\mathbf{c}_i\|_1 + \max(\{\|\mathbf{m}_\alpha^{(i)}\|_1 \mid \alpha \in \Upsilon_i(\mathcal{D}) \setminus \{0\}\})$, and \mathbf{x}_i the minimal element in $\tau_i \cap S$ such that $\|\mathbf{x}_i\|_1$ is greater than or equal to n_i . The vector \mathbf{x}_i can be computed as follows: let Q be the non-negative rational solution of the systems of linear equations $\{x_1 + \dots + x_p = n_i, h_{j_1}^{(i)}(x) = 0, \dots, h_{j_{p-1}}^{(i)}(x) = 0\}$ (recall that $h_{j_1}^{(i)}(x) = 0, \dots, h_{j_{p-1}}^{(i)}(x) = 0$ are the equations defining τ_i), then $\mathbf{x}_i = \left\lceil \frac{\|Q\|_1}{\|\mathbf{a}_i\|_1} \right\rceil \mathbf{a}_i$.

By the proof of Theorem 9, $\mathcal{C}_S \setminus S \subset \mathcal{X}$, with $\mathcal{X} = \{\sum_{i \in [q]} \lambda_i \mathbf{x}_i \mid 0 \leq \lambda_i \leq 1\}$. Algorithm 2 shows the process to computed the set of gaps of S . Note that several of its steps can be computed in a parallel way.

We illustrate Algorithm 2 in the following example. Besides, we confirm our handmade computations by using our free software [7].

Algorithm 2: Computing the set of gaps of a \mathcal{C} -semigroup.

Input: The minimal generating set Λ_S of a \mathcal{C} -semigroup $S \subset \mathbb{N}^p$.

Output: Set of gaps of S .

```

begin
   $\mathcal{H} \leftarrow \emptyset$ ;
   $q \leftarrow$  number of extremal rays of  $L(S)$ ;
  for all the  $i \in [q]$  do
     $\mathbf{c}_i \leftarrow$  conductor of  $\tau_i \cap S$ ;
   $\mathcal{D} \leftarrow \{\sum_{i \in [q]} \lambda_i \mathbf{s}_i \mid 0 \leq \lambda_i \leq 1\} \cap \mathbb{N}^p$ ;
  for all the  $i \in [q]$  do
     $\Upsilon = \{\alpha_1, \dots, \alpha_j\} \leftarrow \Upsilon_i(\mathcal{D}) \setminus \{0\}$ ;
    for all the  $h \in [j]$  do
       $\mathbf{m}_h \leftarrow$  the element in  $v_i(\alpha_h) \cap S$  with minimal 1-norm;
       $n \leftarrow \|\mathbf{c}_i\|_1 + \max(\{\|\mathbf{m}_1\|_1, \dots, \|\mathbf{m}_j\|_1\})$ ;
       $\mathbf{x}_i \leftarrow$  minimal element in  $\tau_i \cap S$  with  $n \leq \|\mathbf{x}_i\|_1$ ;
     $\mathcal{X} \leftarrow \{\sum_{i \in [q]} \lambda_i \mathbf{x}_i \mid 0 \leq \lambda_i \leq 1\} \cap \mathbb{N}^p$ ;
    while  $\mathcal{X} \neq \emptyset$  do
       $\mathcal{Q} \leftarrow \text{First}(\mathcal{X})$ ;
      if  $\mathcal{Q} \notin S$  then
         $\mathcal{H} \leftarrow \mathcal{H} \cup \{\mathcal{Q}\}$ ;
       $\mathcal{X} \leftarrow \mathcal{X} \setminus \{\mathcal{Q}\}$ ;
    return  $\mathcal{H}$  set of gaps of  $S$ .
  
```

Example 12 Consider the \mathcal{C}_S -semigroup S defined in example 10. So, $\mathbf{s}_1 = \mathbf{c}_1 = (2, 0, 0)$, $\mathbf{s}_2 = \mathbf{c}_2 = (4, 2, 4)$, $\mathbf{s}_3 = (0, 1, 0)$ and $\mathbf{c}_3 = (0, 0, 0)$. The set \mathcal{D} is

- $(0, 0, 0), (0, 1, 0), (1, 0, 0), (1, 1, 0), (1, 1, 1), (2, 0, 0), (2, 1, 0), (2, 1, 1),$
- $(2, 1, 2), (2, 2, 2), (3, 1, 1), (3, 1, 2), (3, 2, 2), (3, 2, 3), (4, 1, 2), (4, 2, 2),$
- $(4, 2, 3), (4, 2, 4), (4, 3, 4), (5, 2, 3), (5, 2, 4), (5, 3, 4), (6, 2, 4), (6, 3, 4)\}.$

For example, for the extremal ray τ_1 , $\Upsilon_1(\mathcal{D})$ is the set

$$\{(0, 0), (0, 2), (0, 4), (1, 1), (1, 3), (2, 0), (2, 2), (2, 4)\},$$

and $\cup_{\alpha \in \Upsilon_1(\mathcal{D}) \setminus \{0\}} \{\mathbf{m}_\alpha^{(1)}\}$ is

$$\{(0, 1, 0), (3, 1, 1), (3, 1, 2), (3, 2, 2), (3, 2, 3), (4, 2, 4), (4, 3, 4)\}$$

For τ_2 and τ_3 ,

$$\cup_{\alpha \in \Upsilon_2(\mathcal{D}) \setminus \{0\}} \{\mathbf{m}_\alpha^{(2)}\} = \{(0, 1, 0), (3, 1, 2), (1, 1, 0), (3, 2, 3), (2, 0, 0), (2, 1, 0), (6, 3, 5), (3, 1, 1)\}$$

$$\cup_{\alpha \in \Upsilon_3(\mathcal{D}) \setminus \{0\}} \{\mathbf{m}_\alpha^{(3)}\} = \{(1, 1, 0), (1, 2, 1), (2, 0, 0), (2, 3, 1), (2, 4, 2), (3, 1, 1), (3, 1, 2), (3, 2, 3), (4, 2, 2), (4, 3, 3), (4, 2, 4), (5, 3, 4), (6, 2, 4)\}$$

Then $n_1 = 13$, $n_2 = 24$ and $n_3 = 12$, and $\mathbf{x}_1 = (14, 0, 0)$, $\mathbf{x}_2 = (10, 5, 10)$ and $\mathbf{x}_3 = (0, 13, 0)$. Therefore, the set of gaps of S is,

$$\{(1, 0, 0), (1, 1, 1), (2, 1, 1), (2, 1, 2), (2, 2, 1), (2, 2, 2), (2, 3, 2), (4, 1, 2), (4, 2, 3), (5, 2, 4), (5, 3, 5), (8, 4, 7)\}.$$

By using our implementation of Algorithm 2, we obtain the same gaps:

```
In [1]: ComputeGaps ([[2, 0, 0], [4, 2, 4], [0, 1, 0], [3, 0, 0],
                    [6, 3, 6], [3, 1, 1], [4, 1, 1], [3, 1, 2], [1, 1, 0], [3, 2, 3],
                    [1, 2, 1]])
Out[1]: [[1, 0, 0], [1, 1, 1], [2, 1, 1], [2, 1, 2],
         [2, 2, 1], [2, 2, 2], [2, 3, 2], [4, 1, 2], [4, 2, 3],
         [5, 2, 4], [5, 3, 5], [8, 4, 7]]
```

5 Embedding dimension of \mathcal{C} -semigroups

In [8], it is proved that the embedding dimension of an \mathbb{N}^p -semigroup is greater than or equal to $2p$, and this bound holds. Furthermore, a conjecture about a lower bound of embedding dimension of any \mathcal{C} -semigroup is proposed. In this section, we determine a lower bound of the embedding dimension $e(S)$ of a given \mathcal{C} -semigroup S by studying its elements belonging to \mathcal{A} .

As in previous sections, let $\mathcal{C} \subset \mathbb{N}^p$ be a finitely generated cone and τ_1, \dots, τ_q its extremal rays. For any $i \in [q]$, \mathbf{a}_i is the generator of $\tau_i \cap \mathbb{N}^p$, \mathcal{A} is the finite set $\{\sum_{i \in [q]} \lambda_i \mathbf{a}_i \mid 0 \leq \lambda_i \leq 1\} \cap \mathbb{N}^p$ and Γ_i denotes the minimal generating set of the semigroup $T_i \subset \mathbb{N}^{p-1}$ generated by $\Upsilon_i(\mathcal{A})$. Given a \mathcal{C} -semigroup S , consider $\Lambda'_S := \{\mathbf{s}_{t_1}, \dots, \mathbf{s}_{t_k}\}$ the set of minimal generators of S belonging to $\mathcal{A} \setminus \cup_{i \in [q]} \tau_i$, and $M_l := \{i \in [q] \mid \Upsilon_i(\mathbf{s}_{t_l}) \in \Gamma_i\}$ for $l \in [k]$.

The following result provides us with a lower bound for the embedding dimension of any \mathcal{C} -semigroup.

Proposition 13 *Given a \mathcal{C} -semigroup $S \subset \mathbb{N}^p$, then*

$$e(S) \geq \sum_{i \in [q]} (e(S_i) + e(T_i)) + k - \sum_{i \in [k]} \#(M_i). \tag{2}$$

Proof From Theorem 9, for any $i \in [q]$, there exist $e(S_i)$ minimal generators of S in τ_i . Moreover, for each element $\gamma \in \Gamma_i$, there is at least an element of Λ_S in $v_i(\gamma)$. Note that, for every $\mathbf{s} \in \Lambda_S \setminus \mathcal{A}$, there is no $\gamma \in \Gamma_i$ and $\gamma' \in \Gamma_j$ such that $\mathbf{s} \in v_i(\gamma) \cap v_j(\gamma')$, since for any $i, j \in [q]$, $\gamma \in \Gamma_i$ and $\gamma' \in \Gamma_j$, the intersection $v_i(\gamma) \cap v_j(\gamma')$ is empty or belongs to \mathcal{A} . However, if $\mathbf{s} \in \Lambda'_S$, then it is possible that \mathbf{s} belongs to two (or more) different lines $v_i(\gamma)$ and $v_j(\gamma')$ with $\gamma \in \Gamma_i$ and $\gamma' \in \Gamma_j$ (in that case, $v_i(\gamma) \cap v_j(\gamma') = \{\mathbf{s}\}$). Thus, the value of $\#(M_l)$ indicates the number of different lines $v_i(\gamma_i)$ with $\gamma_i \in \Gamma_i$ to which $\mathbf{s}_{t_l} \in \Lambda'_S$ belongs. So, counting the minimal amount of elements needed to have at least one minimal generator in each

line $v_i(\gamma)$ for each $\gamma \in \Gamma_i$ and $i \in [q]$, we have that the embedding dimension of S is greater than or equal to $\sum_{i \in [q]} (e(S_i) + e(T_i)) + k - \sum_{i \in [k]} \#(M_i)$. \square

Example 14 Consider the \mathcal{C}_S -semigroup S given in example 10. In that case, $\Lambda'_S = \{(3, 1, 2), (1, 1, 0)\}$, $\#(M_1) = 2$ (i.e. $\Upsilon_i(3, 1, 2) \in \Gamma_i$ for $i = 1, 2$), and $\#(M_2) = 2$ ($\Upsilon_1(1, 1, 0) \in \Gamma_1$ and $\Upsilon_2(1, 1, 0) \in \Gamma_3$). So, $\sum_{i \in [q]} (e(S_i) + e(T_i)) + k - \sum_{i \in [k]} \#(M_i) = 5 + 7 + 2 - 2 - 2 = 10$ that is smaller than $e(S) = 11$.

Given any bound, the first interesting question about it is if the bound is reached for some \mathcal{C} -semigroup. The answer is affirmative for (2), and this fact is formulated as follows.

Lemma 15 Let $\mathcal{C} \subset \mathbb{N}^p$ be an integer cone generated by $\{\mathbf{a}_1, \dots, \mathbf{a}_q\}$ and let S_1, \dots, S_q be the non proper numerical semigroups minimally generated by $\{n_1^{(i)}, \dots, n_{e(S_i)}^{(i)}\}$ for each $i \in [q]$. Let $\Lambda'' \subset \mathcal{C} \setminus \cup_{i \in [q]} \tau_i$ be a finite set satisfying:

- for every $\gamma \in \Gamma_i$ and $i \in [q]$, there exists a unique $\mathbf{d} \in \Lambda''$ such that $\Upsilon_i(\mathbf{d}) = \gamma$,
- for every $\mathbf{d} \in \Lambda''$, $\Upsilon_i(\mathbf{d}) \in \Gamma_i$ for some $i \in [q]$.

Then, the embedding dimension of the \mathcal{C} -semigroup S generated by

$$\Lambda'' \cup \bigcup_{i \in [q]} \{n_1^{(i)} \mathbf{a}_i, \dots, n_{e(S_i)}^{(i)} \mathbf{a}_i\}$$

is

$$\sum_{i \in [q]} (e(S_i) + e(T_i)) + k - \sum_{i \in [k]} \#(M_i),$$

where k is the cardinality of $\Lambda'_S = \{\mathbf{s}_1, \dots, \mathbf{s}_k\}$, the set of minimal generators of S belonging to $\mathcal{A} \setminus \cup_{i \in [q]} \tau_i$, and $M_l = \{i \in [q] \mid \Upsilon_i(\mathbf{s}_l) \in \Gamma_i\}$ for $l \in [k]$.

Proof By the hypothesis, there are exactly $\sum_{i \in [q]} e(T_i) + k - \sum_{i \in [k]} \#(M_i)$ minimal generators of S outside its extremal rays, and $\sum_{i \in [q]} e(S_i)$ belonging to its extremal rays. \square

Example 16 Let $S \subset \mathbb{N}^3$ be the semigroup minimally generated by

$$\Lambda_S = \{(2, 0, 0), (4, 2, 4), (0, 2, 0), (3, 0, 0), (6, 3, 6), (0, 3, 0), (3, 1, 1), (3, 1, 2), (1, 1, 0), (3, 2, 3), (1, 2, 1)\}.$$

Note that the cone \mathcal{C}_S is the same as the cone in example 10. So, $\mathcal{A}, \Gamma_1, \Gamma_2$ and Γ_3 are the sets given in that example. For the semigroup S , $\Upsilon_1(\{(3, 1, 1), (3, 1, 2), (1, 1, 0)\}) = \Gamma_1$, $\Upsilon_2(\{(3, 1, 2), (3, 2, 3)\}) = \Gamma_2$ and $\Upsilon_3(\{(1, 1, 0), (1, 2, 1)\}) = \Gamma_3$. Since $(1, 1, 0), (3, 1, 2) \in \mathcal{A}$, $e(S) = 11 = 6 + 7 + 2 - 2 - 2 = \sum_{i \in [3]} (e(S_i) + e(T_i)) + 2 - \sum_{i \in [2]} \#(M_i)$.

Fix a cone \mathcal{C} , studying the different possibilities to select sets of points $K \subset \mathcal{C}$ such that $\cup_{i \in [q]} \Gamma_i$ is the union of the minimal generating set of the semigroup given by $\cup_{Q \in K} \Upsilon_i(Q)$ (for i from 1 to q), we can state results like the following:

Corollary 17 *Let S_1, \dots, S_q be the non proper numerical semigroups minimally generated by $\{n_1^{(i)}, \dots, n_{e(S_i)}^{(i)}\}$ for each $i \in [q]$, and $\Lambda'' \subset \mathcal{C}$ satisfying the hypothesis of Lemma 15. Thus, if $\Lambda'' \subset \mathcal{C} \setminus \mathcal{A}$, then the embedding dimension of the \mathcal{C} -semigroup generated by $\Lambda'' \cup \cup_{i \in [q]} \{n_1^{(i)} \mathbf{a}_i, \dots, n_{e(S_i)}^{(i)} \mathbf{a}_i\}$ is $\sum_{i \in [q]} (e(S_i) + e(T_i))$.*

Finally, we illustrate the above result with an example.

Example 18 Let $S \subset \mathbb{N}^3$ be the semigroup minimally generated by

$$\Lambda_S = \{(2, 0, 0), (4, 2, 4), (0, 2, 0), (3, 0, 0), (6, 3, 6), (0, 3, 0), (3, 1, 1), (4, 1, 2), (5, 2, 4), (2, 1, 0), (1, 2, 0), (3, 2, 3), (1, 2, 1)\}.$$

Again, the cone \mathcal{C}_S is the cone appearing in example 10. Note that the elements $(2, 0, 0)$ and $(3, 0, 0)$ are in S_1 , $(4, 2, 4)$ and $(6, 3, 6)$ belong to S_2 , and $(0, 2, 0)$ and $(0, 3, 0)$ are in S_3 . Moreover, $\Upsilon_1(\{(3, 1, 1), (4, 1, 2), (2, 1, 0)\}) = \Gamma_1$, $\Upsilon_2(\{(5, 2, 4), (3, 2, 3)\}) = \Gamma_2$, $\Upsilon_3(\{(1, 2, 0), (1, 2, 1)\}) = \Gamma_3$, and $\Lambda_S \setminus \cup_{i \in [q]} \tau_i \subset \mathcal{C}_S \setminus \mathcal{A}$. As previous corollary asserts, $e(S) = 13 = 6 + 7 = \sum_{i \in [3]} (e(S_i) + e(T_i))$.

6 On the decomposition of a \mathcal{C} -semigroup in terms of irreducible \mathcal{C} -semigroups

We define the set of pseudo-Frobenius of a \mathcal{C} -semigroup S as $\text{PF}(S) = \{\mathbf{a} \in \mathcal{H}(S) \mid \mathbf{a} + (S \setminus \{0\}) \subset S\}$ (recall that $\mathcal{H}(S) = \mathcal{C} \setminus S$), and the set of special gaps of S as $\text{SG}(S) = \{\mathbf{a} \in \text{PF}(S) \mid 2\mathbf{a} \in S\}$. Note that the elements \mathbf{a} of $\text{SG}(S)$ are those elements in $\mathcal{C} \setminus S$ such that $S \cup \{\mathbf{a}\}$ is again a \mathcal{C} -semigroup.

A \mathcal{C} -semigroup is \mathcal{C} -reducible (simplifying reducible) if it can be expressed as an intersection of two \mathcal{C} -semigroups containing it properly (see [11]). Equivalently, S is \mathcal{C} -irreducible (simplifying irreducible) if and only if $|\text{SG}(S)| \leq 1$. A decomposition of a \mathcal{C} -semigroup S in terms of irreducible \mathcal{C} -semigroups is to express S as intersection of irreducible \mathcal{C} -semigroups. This definition generalizes the definitions of irreducible numerical semigroups (see [13]) and irreducible \mathbb{N}^p -semigroups (see [2]).

Our decomposition method into irreducible is based on adding to a \mathcal{C} -semigroup elements of $\text{SG}(S)$. If we repeat this operation, we always reach an irreducible \mathcal{C} -semigroup or the cone \mathcal{C} . Since the set of gaps $\mathcal{H}(S)$ is finite, this process can be performed only a finite number of times. This allows us to state the following algorithm inspired by [13, Algorithm 4.49].

By definition, the set $\text{SG}(S)$ is obtained from $\text{PF}(S)$. If S is determined by its minimal generating set, then $\text{PF}(S)$ can be computed from the set $\mathcal{H}(S)$ obtained with Algorithm 2, or using the two different ways given in [11, Corollary 9 and Example 10].

Example 19 Consider the \mathcal{C} -semigroup S given in examples 10 and 12. It is minimally generated by

Algorithm 3: Computing a decomposition into \mathcal{C} -semigroups.

Input: The minimal generating set Λ_S of a \mathcal{C} -semigroup $S \subset \mathbb{N}^p$.

Output: A decomposition of S into irreducible \mathcal{C} -semigroups.

```

begin
  I ← ∅;
  C ← {S};
  while C ≠ ∅ do
    B ← {S' ∪ {a} | S' ∈ C, a ∈ SG(S')};
    B ← B \ {S' ∈ B | ∃S̄ ∈ I with S̄ ⊂ S'};
    I ← I ∪ {S' ∈ B | S' is irreducible};
    C ← {S' ∈ B | S' reducible};
  return I.

```

$$\Lambda_S = \{(2, 0, 0), (4, 2, 4), (0, 1, 0), (3, 0, 0), (6, 3, 6), (3, 1, 1), (4, 1, 1), (3, 1, 2), (1, 1, 0), (3, 2, 3), (1, 2, 1)\},$$

with

$$\mathcal{H}(S) = \{(1, 0, 0), (1, 1, 1), (2, 1, 1), (2, 1, 2), (2, 2, 1), (2, 2, 2), (2, 3, 2), (4, 1, 2), (4, 2, 3), (5, 2, 4), (5, 3, 5), (8, 4, 7)\}.$$

Hence, $\text{PF}(S) = \{(2, 2, 1), (2, 3, 2), (4, 1, 2), (8, 4, 7)\}$, and $\text{SG}(S)$ is equal to $\text{PF}(S)$.

Applying Algorithm 3 to S , we obtain the decomposition into six irreducible \mathcal{C} -semigroups, $S = S_1 \cap \dots \cap S_6$ where

- $S_1 = \langle (3, 0, 0), (2, 0, 0), (1, 1, 0), (0, 1, 0), (4, 1, 1), (3, 1, 1), (3, 1, 2), (4, 1, 2), (1, 2, 1), (2, 2, 1), (2, 2, 2), (3, 2, 3), (4, 2, 4), (6, 3, 6) \rangle$;
- $S_2 = \langle (3, 0, 0), (2, 0, 0), (1, 1, 0), (0, 1, 0), (4, 1, 1), (3, 1, 1), (2, 1, 2), (3, 1, 2), (1, 2, 1), (2, 2, 1), (3, 2, 3) \rangle$;
- $S_3 = \langle (1, 0, 0), (0, 1, 0), (2, 1, 1), (3, 1, 2), (1, 2, 1), (3, 2, 3), (4, 2, 4), (5, 3, 5), (6, 3, 6) \rangle$;
- $S_4 = \langle (3, 0, 0), (2, 0, 0), (1, 1, 0), (0, 1, 0), (2, 1, 1), (1, 1, 1), (3, 1, 2), (4, 1, 2), (3, 2, 3), (4, 2, 4), (6, 3, 6) \rangle$;
- $S_5 = \langle (3, 0, 0), (2, 0, 0), (1, 1, 0), (0, 1, 0), (2, 1, 1), (1, 1, 1), (3, 1, 2), (3, 2, 3), (4, 2, 4), (5, 2, 4), (6, 3, 6) \rangle$;
- $S_6 = \langle (3, 0, 0), (2, 0, 0), (1, 1, 0), (0, 1, 0), (4, 1, 1), (3, 1, 1), (2, 1, 2), (3, 1, 2), (1, 2, 1), (3, 2, 3), (4, 2, 3) \rangle$;

To get these semigroups we have used our implementation in [7] by typing the following

```

Csemigroup([ [2, 0, 0], [4, 2, 4], [0, 1, 0], [3, 0, 0], [6, 3, 6],
[3, 1, 1], [4, 1, 1], [3, 1, 2], [1, 1, 0], [3, 2, 3], [1, 2, 1] ]) .
DecomposeIrreducible()

```

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Declarations

Conflict of interests The authors declare that there is no conflict of interests regarding the publication of this paper.

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