

MASTER

Realization of a real-time 16 kbit/s speech coder and decoder on a single digital signal processor chip each

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DEPARTMENT OF ELECTRICAL ENGINEERING
EINDHOVEN UNIVERSITY OF TECHNOLOGY
TELECOMMUNICATIONS DIVISION EC

REALIZATION OF A REAL-TIME 16 KBIT/S
SPEECH CODER AND DECODER ON A SINGLE
DIGITAL SIGNAL PROCESSOR CHIP EACH

by R.P.J. Kleuters

Report of the graduation work
accomplished from 8-9-1986 to 17-7-1987.
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The department of electrical engineering of the Eindhoven
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SUMMARY

This report describes the design, implementation and testing of a real-time 16 kbit/s speech coder and decoder. The algorithm used is simple enough for the coder and decoder to be implemented each on a single digital signal processor chip (TMS32010).

In the transmitter the up to 4 kHz bandlimited speech signal is first split into eight mutually exclusive subband signals, each having a spectral bandwidth of 500 Hz. The sampling rates of the subband signals are reduced from 8 kHz to 1 kHz (being the Nyquist rate). For this subband splitting and sampling rate reduction, a very efficient technique called Quadrature Mirror Filtering is used. Next, each subband signal is independently coded according to its perceptual contribution to the overall subjective quality. For coding the subband signals, Pulse Code Modulation, with backward adaption of the quantization stepsize, is employed. Furthermore for each subband signal, the number of code bits is determined by a semi-adaptive bit-allocation algorithm. Finally, the coded subband samples together with the necessary side information and synchronization bits are multiplexed into a 16 kbit/s serial data stream.

In the receiver, the inverse operations (demultiplexing, decoding and speech signal reconstruction) take place.

The quality performance of the implemented Subband Coder, in particular the intelligibility of the reconstructed speech, is acceptable for our purposes. There is only little difference with unprocessed up to 3 kHz bandlimited speech. Furthermore, the capacities of each digital signal processor are nearly completely utilized by the implemented Subband Coder and Decoder algorithm, which means that an appropriate type of digital signal processor has been chosen.

LIST OF ABBREVIATIONS AND SYMBOLS

AIB	Analog Interface Board; part of Texas Instruments development system.
A/D	Analog to Digital (conversion).
AQB	Adaptive Quantization Backward.
CRT	Cathode Ray Tube.
D/A	Digital to Analog (conversion).
d_c	Stepsize adaption control factor.
$d(n)$	$\log \Delta(n)$.
$\Delta(n)$	Stepsize for (de-)quantization of a subband sample.
DPCM	Differential PCM.
DSP	Digital Signal Processor.
EVM	Evaluation Module; part of Texas Instruments development system.
F	Number of bits in a frame of the multiplexed signal.
FAW	Frame Alignment Word (in multiplexing operations).
FIR	Finite Impulse Response (filter).
f_s	Sampling rate.
γ	Stepsize adaption leakage factor.
$I(n)$	Two's complement B-bit code word; coded subband sample.
I/O	Input/Output.
m	Number of bits in a FAW.
$M(.)$	Stepsize adaption parameter.
$m(.)$	$\log M(.)$.
p	Bit error rate.
PC	Personal Computer.
PCM	Pulse Code Modulation.
PDF	Probability Density Function.
PROM	Programmable Read Only Memory.
QMF	Quadrature Mirror Filter.
Qx	Fixed-point format in which a number is represented in the TMS32010. A number is represented with 1 sign bit, 15-x integer bits and x fractional bits.

RAM Random Access Memory.
rms Root mean square.
SBC Subband Coder.
s(n) 8 kHz digital speech input to transmitter.
ŝ(n) 8 kHz reconstructed digital speech output from
 receiver.
s(t) Analog speech input to transmitter.
ŝ(t) Reconstructed analog speech output from receiver.
t₁ Mean time to acquire frame alignment.
t₂ Mean time between false indications of lost
 frame alignment.
x(n) Quantizer input; uncoded subband sample.
x̂(n) Quantized subband sample.

1. INTRODUCTION

In the Telecommunications Division of the Eindhoven University of Technology work is done on a cooperative project with the University of Dar-es-salaam. This project concerns the distribution of audio-visual information via satellite channels to be used for tele-education, information dissemination and news distribution to the rural population and to isolated institutions in Africa. To reduce costs, services suitable for standard narrowband (64 kbit/s) distribution channels are necessary. Examples are transmission of still pictures plus speech, teletext plus speech, silent-teletext and scribophone. For more information on the development project, see [1].

This report deals with the realization of a part of the above mentioned services, namely a real-time speech coder and decoder. In this project 16 kbit/s of a 64 kbit/s channel is reserved for the transmission of speech. Another design requirement is to keep the costs as low as possible. Considering the enormous advances in digital technology, coupled with the increasing economics of digital hardware, it was decided to implement the coder and decoder in software on one digital signal processor (DSP) chip each. A technique that can allow this and that can realize a fairly good quality of the reconstructed speech at the receive side is known as Subband Coding [2,3,4,5].

As the first work on the Subband Coder (SBC) project a study has been made about the kind of Subband Coding which best meets the requirements [6]. From this work a proposal resulted for the type of Subband Coding to be used. During further work [7,8] implementations of functional blocks of this proposal were designed and tested.

To realize a properly working Subband Coder and Decoder these implementation designs had to be adjusted and extended. To meet the DSP capacities the software also had to be optimized. Besides these alterations of the implementation designs and the software, the hardware communication between the transmitter and the receiver system had to be realized.

In this report the ultimate implementation design and hardware environment of a properly working real-time Subband Coder and Decoder realization are described.

First some general design considerations are discussed (chapter 2). This is followed by a description of the functional blocks of the Subband Coder and Decoder and the work that is done on the implementation design of them (chapter 3 to 6). After that the employed hardware will be considered (chapter 7). Finally the utilization of the transmitter and receiver DSP is discussed (chapter 8), followed by some suggestions for possible performance improvement (chapter 9) and the conclusions (chapter 10).

2. GENERAL DESIGN CONSIDERATIONS

2.1. Introduction

In this chapter some general aspects concerning the design of a Subband Coder are considered. First we will discuss the principle of Subband Coding. Then some features of the used digital signal processor are showed. This is followed by a description of the development equipment, and finally the design strategy is outlined.

2.2. General principle of Subband Coding

A general block diagram of a Subband Coder and Decoder is shown in Fig. 1.

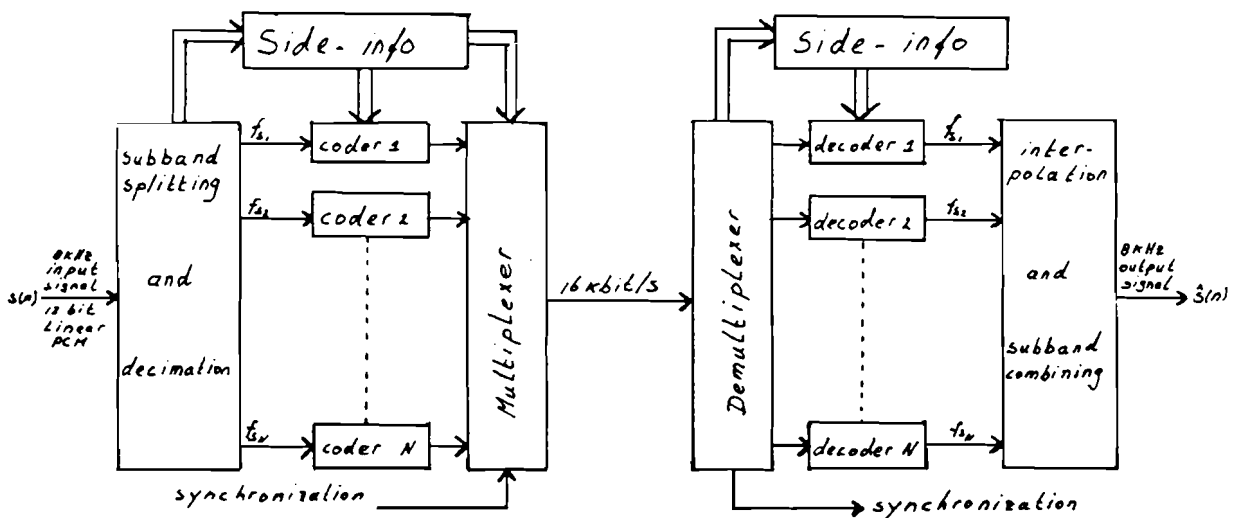


Fig. 1: General block diagram of a Subband Coder and Decoder.

The digital input to the transmitter, $s(n)$, is obtained by sampling a bandlimited (up to 4 kHz) analog speech signal at a rate of 8 kHz, followed by 12 bit linear PCM analog-digital conversion. This signal is filtered into a number of N subband signals, each with a different spectrum that is part of the baseband frequency range 0-4 kHz. After the bandsplitting the sampling rates, f_{s_i} , of these subband signals can be reduced (decimation). This is possible because

the spectrum of each subband signal is narrower than the spectrum of the input signal. The minimum sampling frequency is twice the bandwidth of a subband signal. Decimation by an integer factor M can be achieved by retaining only one out of every M filter output samples. Then the subband signals are independently quantized and PCM coded into a number of bits that is determined by the bit-allocation used. In any case the bit-allocation is so as to allow an ultimate bit rate of 16 kbit/s. Finally the different subband signals together with the necessary side and synchronization information will be multiplexed into a serial 16 kbit/s data stream, representing the output signal of the transmitter. In the receiver the inverse actions will take place. First the incoming data stream is demultiplexed to recover the coded subband signals plus the side and synchronization information. Next the subband signals are decoded in the PCM decoders. The decoding takes place according to the bit-allocation used, which in an adaptive case can be extracted from the received side information. By inserting zero-valued samples between the subband samples, followed by a filtering identical to that in the transmitter, the sampling rate in the subband signals will be restored to the original sampling rate of 8 kHz. During the filtering interpolation takes place, which gives the inserted zero's an appropriate value. Increasing the sampling rate with an integer factor M is accomplished by filling in $M-1$ zero-valued samples between each pair of filter input samples. By combining these final obtained subband signals the original digital speech signal, $S(n)$, can be reconstructed.

The main advantages of Subband Coding relative to other coding techniques are:

- A low bit rate.
- The quantization noise generated within each subband remains confined to that channel and is independent of noise produced in other bands. In this way high level, low frequency quantization noise does not mask low level, high

- frequency sounds, and vice versa.
- Each subband signal may be coded independently according to the perceptual contribution that each band makes to the overall subjective quality. This together with the sampling rate reduction in the subbands is the explanation for the low bit rate that can be achieved.
 - By proper design a real-time implementation of the whole SBC principle is possible on two DSP's; one for the transmitter and one for the receiver.

These are enough reasons for using the Subband Coding technique for the realization of a 16 kbit/s speech coder and decoder that together deliver fairly good quality speech at a fairly low complexity (c.q. costs).

2.3. Capacities of the used DSP

The DSP used for the implementation of the SBC algorithm, is the TMS32010 from Texas Instruments. This TMS32010 is very suitable for speech processing applications. For an extensive documentation of the TMS32010, see [9]. As both the coder and decoder had to be implemented on one DSP chip each, it is not strange that the features and capacities of the TMS32010 have had a great influence on the design work. The most important ones will be discussed below.

2.3.1. Processor speed

The duration of one instruction cycle is 200 ns. Subband Coding is a real-time speech coding technique that in our case every 0.125 ms uses one input sample in the transmitter and creates one output sample in the receiver. So in the transmitter as well as in the receiver the whole coder resp. decoder program has to be passed through in 0.125 ms. So the longest cycle in either program may not count more than 625 instruction cycles.

2.3.2. Program memory and data memory

Program memory consists of up to 4K words of 16-bit width. In our case this concerns off-chip RAM. The available program memory is large enough and has been no constraint on the design work.

On the other hand data memory consists of 144 words of 16-bit width of RAM present on-chip. So data RAM is far too small to contain all variables for the coder or decoder. This problem can be solved by storing some data operands in off-chip program memory and then read them into on-chip data memory when needed. When they are not needed anymore they can be written back to program RAM if necessary. However this solution will slow down execution as reading from and writing to program memory each take 600 ns, whilst most instructions only need 200 ns. So it is necessary to carefully determine which variables are directly stored in data memory and which are stored in program memory. Variables that are not often used and tables for example can best be stored in program memory. For more information, see [7, pp. 39-40]. The available data memory has appeared to be a great constraint in the design of the SBC algorithm.

2.3.3. Arithmetical operations

Addition and subtraction are standard tools for the TMS32010 and will cost only one instruction cycle. Also multiplication can take place in 200 ns due to the presence of a 16*16-bit parallel multiplier.

However, the TMS32010 does not have an explicit divide instruction. A division therefore will have to be broken down into a series of subtracts and shifts with the consequence of a long execution time. So divisions have to be avoided as far as possible.

2.3.4. Fixed-point arithmetic

Computation on the TMS32010 is based on a fixed-point two's complement representation of numbers. Each 16-bit number is evaluated with a sign bit, i integer bits, and 15 minus i fractional bits. Thus the number :

0000 0010 1010 0000
└──────────┬── decimal point

has a value of 2.625. This particular number is said to be represented in a Q8 format (8 fractional bits). Its range is between -128 (1000 0000 0000 0000) and +127.996 (0111 1111 1111 1111). The fractional accuracy of a Q8 number is about 0.004 (one part in $2^{*}8$ or 256).

To reduce quantization noise on the one side and to avoid overflows on the other side, it is very important to select an appropriate representation (i.e. Q-format) for each variable and constant.

On the design work we have started from the principle that the transmitter DSP input and the receiver DSP output have been normalized to the range between -1 and +1 Volt, and therefore a Q15 format has been taken for the representation of these two signals.

2.4. Equipment for testing and realizing the software developments

Two development systems as shown in Fig. 2 are available, one for the transmitter DSP and one for the receiver DSP. Such development system consists of two printed circuit boards, called "Evaluation Module" (EVM) and "Analog Interface Board" (AIB), and a host computer, in our case a Personal Computer from IBM.

The basic purpose of the Evaluation Module and the Analog Interface Board used is to enable the user to develop programs for the TMS32010 digital signal processor and to run

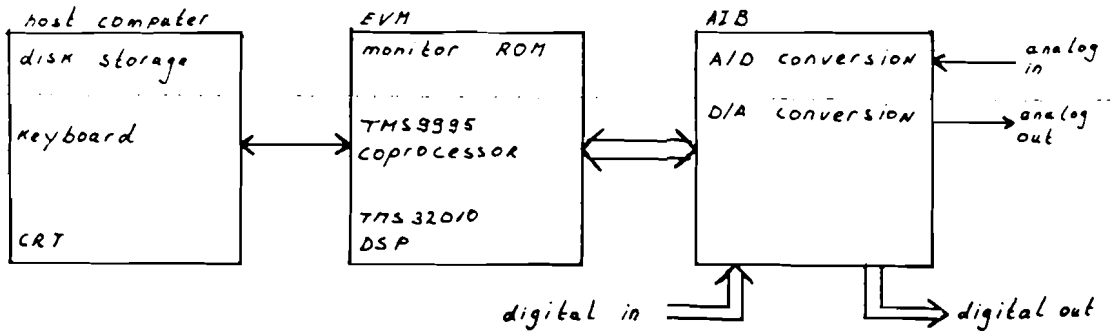


Fig. 2: Development system for the TMS32010.

these programs real-time. The programs can communicate with the analog "outside world" by means of A/D and D/A converters and bandlimiting/interpolation filters resident on the AIB. The program storage, the program design tools and the DSP itself are resident on the EVM board.

The two printed circuit boards of Texas Instruments are controlled by the host computer (IBM PC). This PC can communicate with the EVM via a serial interface. At the IBM PC the serial interface is controlled under BASIC. The EVM board is equipped with two serial interfaces of which only one, called "port 1", will be used. All serial interfaces are bi-directional. The EVM board together with the IBM PC form a program development system with which the user may enter assembler source files, assemble files, run them, execute them with breakpoints and single step programs.

The EVM and the AIB together can form a stand-alone system once these are programmed and started.

The EVM and AIB are well documented in [10] resp. [11]. In [7, pp. 24-32] the possibilities of the development system and the implications of this for a communications program resident in the IBM PC are discussed in more detail.

Later on in the project the development equipment has been extended with a so-called cross-assembler, which allows us to assemble source files on the IBM PC. This cross-assembler is described in [12] and [13].

2.5. Design_strategy

When we look back at the general SBC principle described in 2.2., we can discern the following functional blocks:

- filtering
- bit-allocation
- coding and decoding
- multiplexing and demultiplexing

As a logical consequence the design work, implementation and testing has been done in a modular way.

First the filtering has been realized and tested on proper working. After that the bit-allocation has been added. This was followed by adding the coding and decoding. Finally everything has been supplemented with the multiplexing and demultiplexing.

3. ANALYSIS/SYNTHESIS_FILTER_BANK

3.1. Introduction

An important and critical aspect of the SBC design is the filter bank and its interaction with the sampling rate reduction (decimation) in the transmitter and the sampling rate increase (interpolation) in the receiver. A good reconstruction of the speech signal, $s(n)$, requires a subband splitting in the transmitter and subsequent subband combining in the receiver with an overall frequency response that in magnitude equals 1 as best as possible.

If we leave out of consideration effects such as quantization noise, which are no direct consequence for filtering, there are two effects that can cause the magnitude of the overall frequency response to be not exactly equal to 1 everywhere. Both effects concern the fact that the filters we deal with in practice are not ideal in their frequency response, but contain beside the passband also a transition band and a stop band.

-The reduction of the subband sampling rates is necessary in order to maintain a low (16 kbit/s) bit rate in encoding these signals. This sampling rate reduction introduces aliasing terms [14, pp. 304-305] in each of the subband signals.

-In the reconstruction process the subband signals are combined together. As interpolation in the receiver takes place with a filtering identical to the one in the transmitter, after reconstruction the original signal has twice passed the filter bank represented by the frequency responses $H_1(f)$, $H_2(f)$, ..., $H_N(f)$. So exact reconstruction requires: $|H_1(f)|^2 + |H_2(f)|^2 + \dots + |H_N(f)|^2 \equiv 1$. In places where this requirement is not met, a ripple occurs in the response, called reconstruction ripple. Reconstruction ripple mainly occurs in the transition regions of the filter responses.

Thus concerning the filter bank development, aliasing and reconstruction ripple have to be suppressed or avoided as much as possible to accomplish a good speech signal reconstruction.

In the following the design, implementation and testing of the subband splitting and reconstruction method that has been chosen for our SBC realization is discussed.

3.2. Design considerations

Several methods exist for the splitting of the SBC input signal into a number of mutually exclusive subbands [6]. Some methods [15,16,17] are based on using a parallel bank of bandpass filters, as shown in Fig. 3.

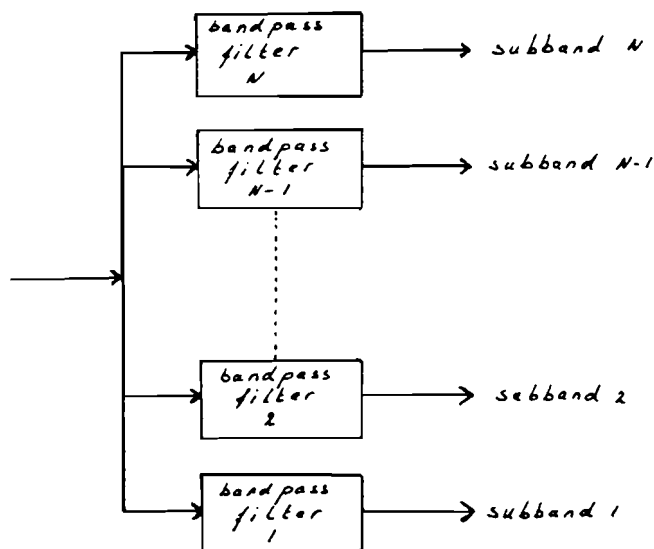


Fig. 3: Parallel bank of bandpass filters for subband splitting.

These filters may be conveniently realized by linear-phase Finite Impulse Response (FIR) networks. However, to meet the design requirements as mentioned in 3.1., very sharp transition band filters, i.e. very high order networks (ca. 200-tap) with a strong stop band attenuation (ca. 45 dB) are needed. Such filters would have a complexity that is far too high, excessive delay and possible performance limitations due to finite precision wordlengths if realized directly.

A more convenient alternative [18,19,20,21] is the use of a tree-structured filter bank. Such a tree configuration works by successively splitting the signal into two subbands at each branching point, using a high-pass/low-pass filter combination. The output from each intermediate filter is downsampled to the appropriate Nyquist rate for the signal and then applied to the next branching level for further spectral division, see Fig. 4.

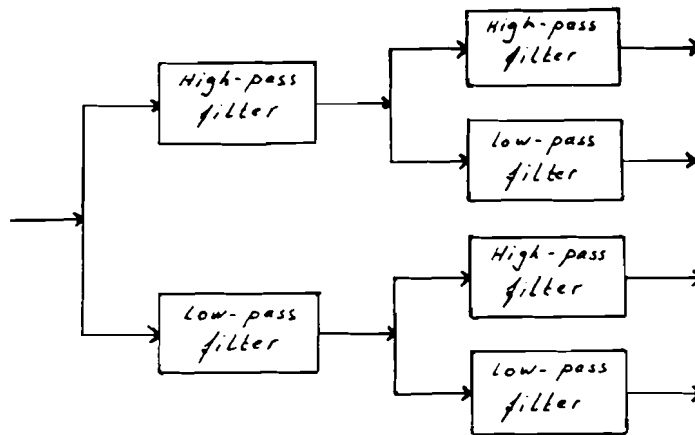


Fig. 4: Tree-structured filter bank for uniformly spaced subbands.

Besides this uniform subband splitting also octave-spaced subbands are possible when using the approach as shown in Fig. 5.

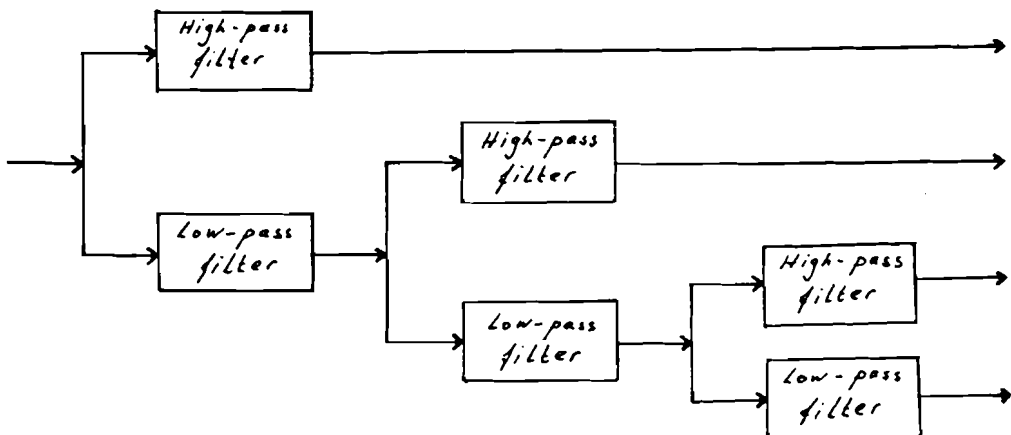


Fig. 5: Tree-structured filter bank for octave-spaced subbands.

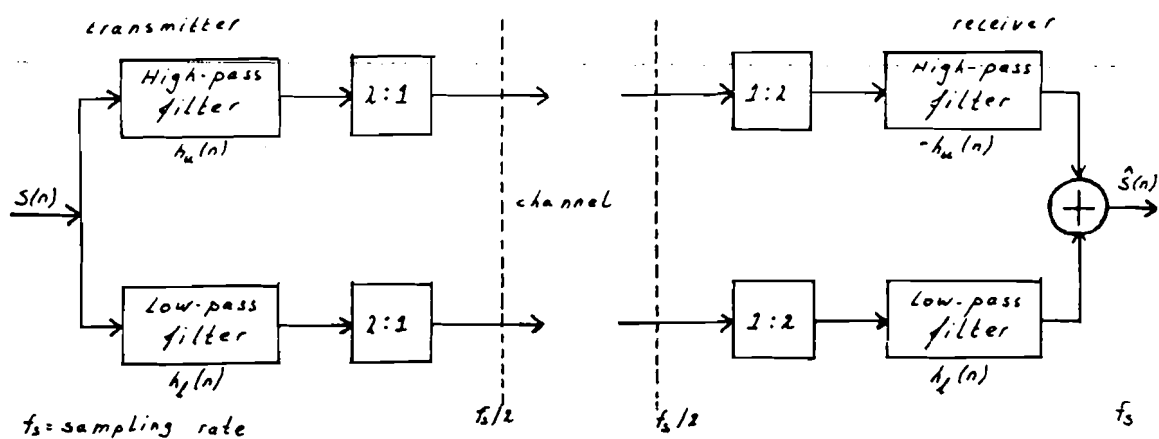
If we would realize the tree-structured filter bank with the usual FIR filters, without taking any special measures, again very high order networks are needed.

However for the bandsplitting a technique called "Quadrature Mirror Filtering" can also be used. Quadrature Mirror Filters (QMF's) have special phase and magnitude characteristics which allow the splitting of a band into two equal width subbands and, upon reconstruction, provide for the cancellation of aliasing effects that occur during downsampling. Furthermore, as each band is splitted into two symmetrical parts, it is not very difficult to see to it that reconstruction ripple is strongly suppressed. Because of these properties, lower order filters (32 to 12 tap), which have fairly wide transition widths, can be used to cover the entire speech band of interest without any spectral gaps in the total response. Therefore, and also for some other reasons to be discussed later, a tree-structured QMF bank has been chosen for the realization of our SBC.

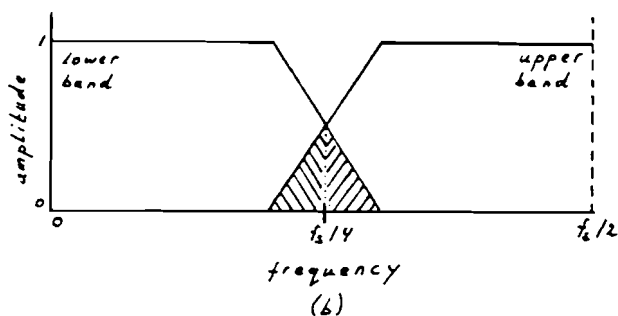
Fig. 6 reviews the basic configuration for a two-band SBC design that will be used for the explanation of the QMF bandsplitting and derivation of its design requirements.

The transmitter input signal $s(n)$ is divided into two equally spaced frequency bands by low-pass and high-pass filters, $h_l(n)$ and $h_u(n)$, respectively. Each subband signal is reduced in sampling rate by a factor of two, i.e. if f_s is the sampling rate of the input signal, $f_s/2$ is the sampling rate of the subband signals. In practice the subband signals are then encoded and multiplexed for transmission, which in this (modular) design stage will be left out of consideration.

In the receiver the subband signals are interpolated back to their original sampling rates with the aid of similar low-pass and high-pass filters. The sum of the two interpolated subband signals, $\hat{s}(n)$, is the reconstructed version of the input signal, $s(n)$.



(a)



(b)

Fig. 6: (a) General block diagram of a two-band SBC (coding and multiplexing left out of consideration)
 (b) Spectral description of the subbands.

To obtain the aliasing cancellation property, the filters $h_l(n)$ and $h_u(n)$ must be symmetrical FIR designs with even numbers of taps, i.e.,

$$h_l(n) = h_u(n) = 0 \quad \text{for } n < 0 \text{ and } n > N-1 \quad (1)$$

where N (even) is the number of taps. The symmetry property implies that

$$h_l(n) = h_l(N-1-n), \quad n=0, 1, 2, \dots, N/2-1, \text{ and} \quad (2a)$$

$$h_u(n) = -h_u(N-1-n), \quad n=0, 1, 2, \dots, N/2-1 \quad (2b)$$

Furthermore it is required that

$$h_u(n) = (-1)^n h_l(n) \quad n=0,1,2, \dots, N-1 \quad (3)$$

which is the mirror image relationship of the filters.

With the above constraints, the aliasing cancellation property of the QMF bank can be verified easily. A derivation is given in the appendix of [20].

To suppress reconstruction ripple as much as possible, the filters $h_l(n)$ and $h_u(n)$ ideally must satisfy also the condition

$$|H_l(f)|^2 + |H_u(f)|^2 \equiv 1 \quad (4)$$

where $H_l(f)$ and $H_u(f)$ are the Fourier transforms of $h_l(n)$ and $h_u(n)$ respectively. This also can be seen from the derivation in the appendix of [20].

The above filter requirement of eq. (4) cannot be met exactly except when $N=2$ and when N approaches infinity. However, it can be very closely approximated for modest values of N . Filter designs which satisfy eq. (2a) and approximate the condition of eq. (4) and the lowpass characteristic can be obtained with the aid of an optimization program. In practice "Hanning window" designs, optimized by the "Hooke and Jeeves algorithm" [6,21] will give satisfactory results. Fig. 7 shows the frequency response characteristics for an $N=32$ -tap filter design, acquired with the above mentioned technique. As can be seen from Fig. 7b, the requirement of eq. (4) is satisfied to within ± 0.025 dB, which is more than satisfactory for good SBC performance.

The QMF technique discussed for the two-band SBC from Fig. 6a can be applied in the same way at each branching point in a SBC that employs a tree-structured filtering into more than two subbands.

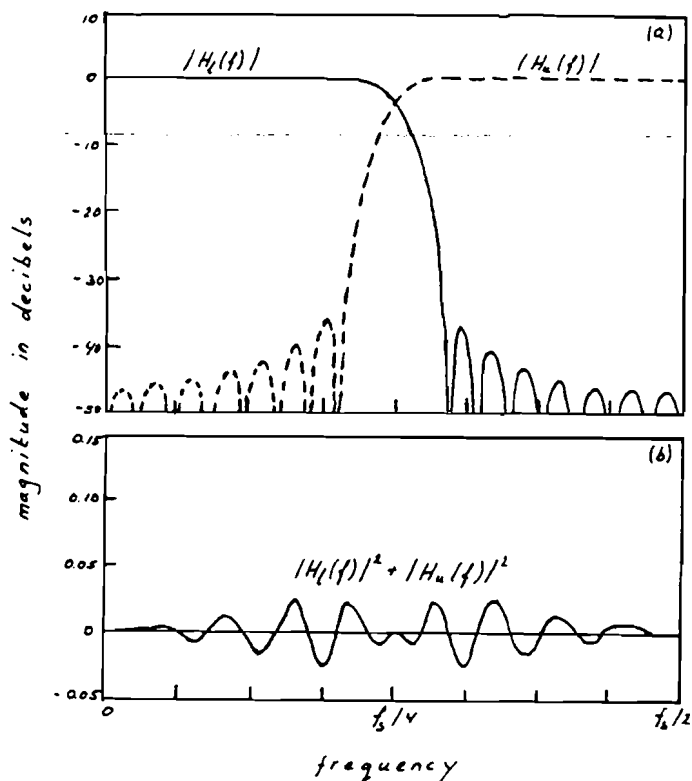


Fig. 7: Frequency response for a 32-tap QMF design [20]
(a) Magnitude responses of the individual filters
(b) Magnitude response of the composite system.

Finally we have to sketch the form of tree-structured QMF bank that has been chosen. As we will see in chapter 4, some adaption in the bit-allocation is necessary to achieve satisfactory reconstruction speech quality. However, if the subbands are unequally spaced, this requires techniques of a very high complexity. Therefore we have chosen uniformly spaced subbands [6,18]. Furthermore, the perceptual quality of the recovered speech will increase as the number of subbands increases. The number of subbands to be used in practice is a trade-off between recovered speech quality, processing complexity and delay in the filter bank. For our purposes an eight-band SBC has been found to represent a good compromise [6].

A sketch of the ultimate filterbank design to be implemented in the transmitter and receiver is shown in Fig. 8.

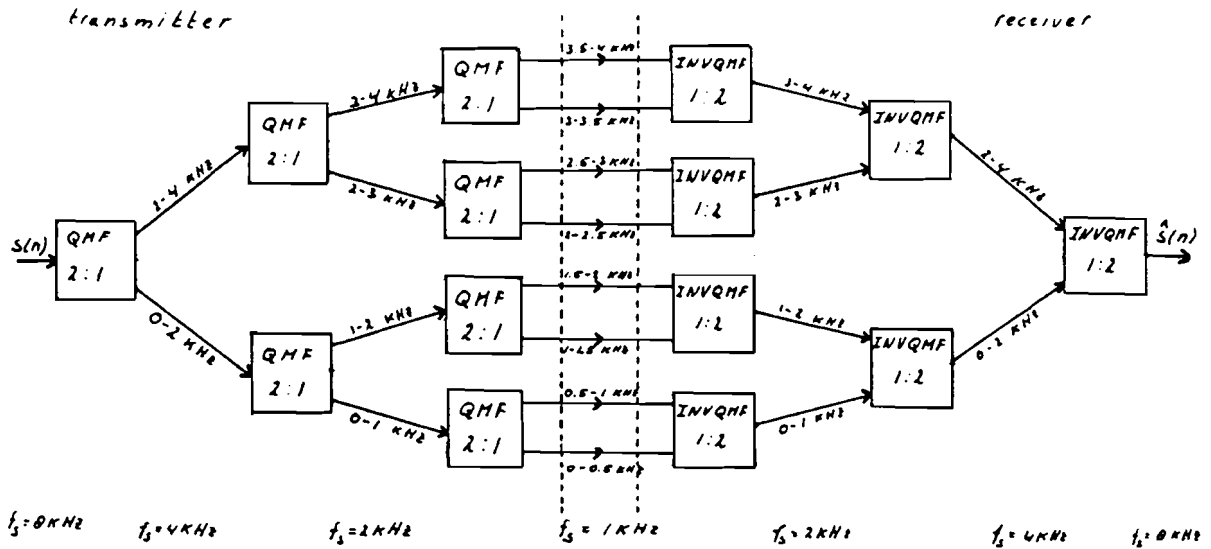


Fig. 8: Eight-Band subband splitting and reconstruction using a tree-structured QMF bank.

3.3. Implementation of the chosen filter bank algorithm

Here again, the basic principle of the implementation of the eight-band tree-structured QMF bank is discussed for the case of a two-band SBC, see Fig. 6a. From the mirror image property of the QMF bank described by eq. (3), we note that the coefficients used for the upper and lower subband filters are identical, except for the signs of the odd numbered coefficients. This property can be used to save a factor of two in computation by sharing the computation between the

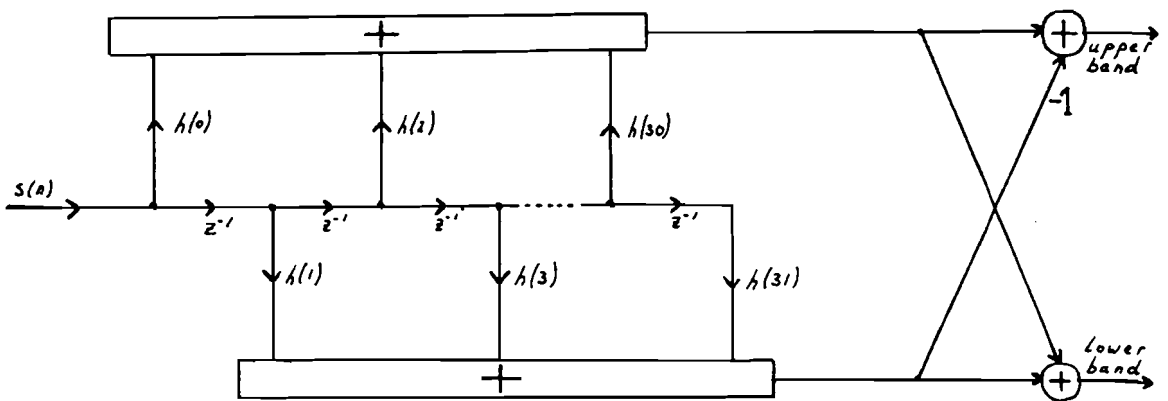


Fig. 9: QMF bank structure that shares computation between upper and lower filters.

filters in the manner described in Fig. 9, where $h_l(n)=h(n)$ and $h_u(n)=(-1)^n h(n)$. The partial sums of products are accumulated separately for the even- and odd-filter coefficient values. The sum of these two partial sums then gives the lower subband signal, and their difference produces the upper subband signal. Since the subband sampling rates are one-half of the input sampling rate, an additional factor of two is gained by computing the sums of products indicated in Fig. 9 once for every other input sample. Thus, each sample is shifted two delays in the shift register of Fig. 9 before being used.

Because of this sampling rate reduction, the filter structure of Fig. 9 can be divided into two parts as shown in Fig. 10.

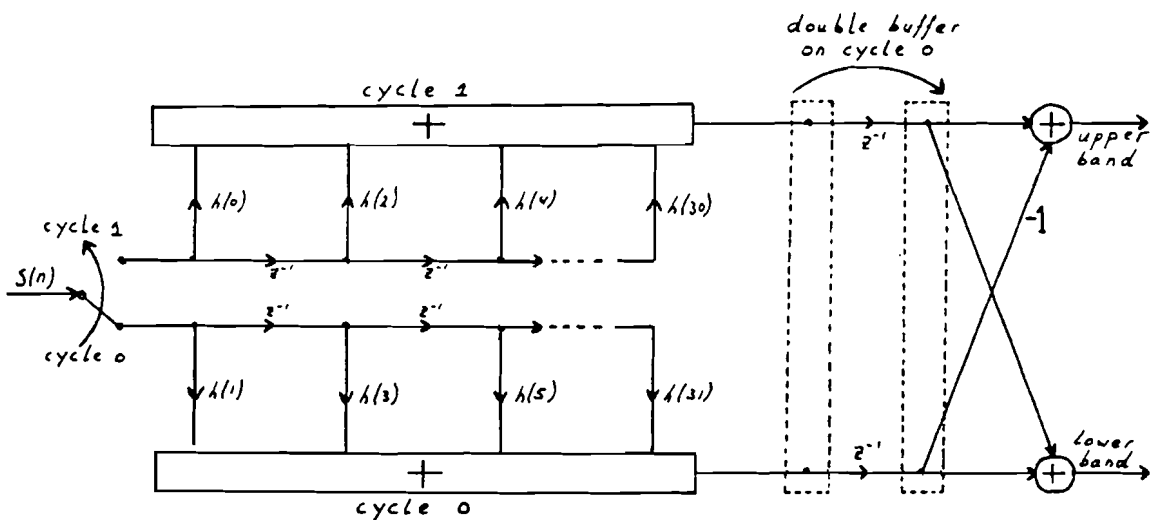


Fig. 10: Polyphase QMF bank structure of a SBC transmitter.

This structure is a two-band version of a more general class of multirate structures sometimes referred to as polyphase structures. As Fig. 10 shows, the input signal is separated into two sets by a commutator. Assuming that the commutator is in the lower position at time $n=-1$, the lower branch receives odd values of $s(n)$, i.e. $s(-1)$, $s(1)$, $s(3)$, ..., and the upper branch receives even values of $s(n)$, i.e. $s(0)$, $s(2)$, $s(4)$, Both branches now operate at one-half of

the original sampling rate. Odd values of $s(n)$ are filtered at odd sample times in the lower branch with an $N/2$ tap filter of odd valued filter coefficients (cycle 0). Similarly, even valued samples of $s(n)$ are filtered in the upper branch with a $N/2$ tap filter of even filter coefficients. Furthermore double buffering is required when data computed in one cycle are needed in another cycle. The outputs of the even and odd filters are computed and stored in the left buffer for cycles 0 and 1. For taking the sum and difference the available data from the right buffer, which have been computed in the previous 0 and 1 cycles, are used. At the beginning of cycle 0 the data are transferred from the left buffer to the right buffer. The sum in the lower branch of Fig. 10 is computed in cycle 0, while the difference is computed in cycle 1.

A similar efficient polyphase structure can be generated for the QMF synthesis bank in the receiver. The resulting structure is shown in Fig. 11, that speaks for itself.

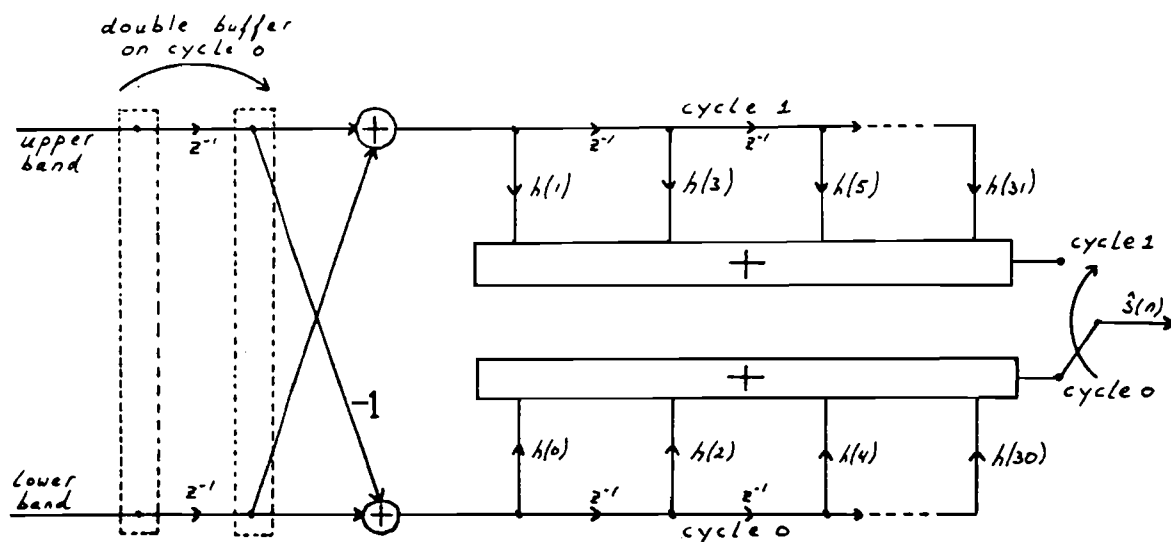


Fig. 11: Polyphase QMF bank structure of a SBC receiver.

So, in this QMF analysis and synthesis implementation, that makes efficient use of the available DSP resources, the computation load is evenly distributed between even and odd time cycles of the input sampling rate. This also has been

one of the reasons for choosing the QMF technique. Table 1 reviews the two-cycle control structure that is used in the implementation of a two-band QMF splitting and reconstruction.

Table 1: Control structure for two-band QMF analysis and synthesis.

Cycle 0	A. Transmitter
	1. Double buffer
	2. Create lower band output by summation
	3. Input one sample of $s(n)$
	4. FIR filter (lower branch)
	B. Receiver
	1. Double buffer
	2. Create lower branch filter input by subtraction
	3. Input one subband sample and store in left buffer
	4. FIR filter (lower branch)
	5. Output one sample of $\hat{s}(n)$

Cycle 1	A. Transmitter
	1. Create upper band output by subtraction
	2. Input one sample of $s(n)$
	3. FIR filter (upper branch)
	B. Receiver
	1. Create upper branch filter input by summation
	2. Input one subband sample and store in left buffer
	3. FIR filter (upper branch)
	4. Output one sample of $\hat{s}(n)$

For the implementation of the eight-band SBC, this technique is repeated at each branching point of the filter tree. As the eight-band SBC is a multirate system with an ultimate sampling rate ratio of eight, it requires an eight-cycle control structure. For each cycle only one out of the eight possible paths in the QMF tree has to be passed through. A modulo-8 counter, called PATH, is used to assign the path in the QMF tree to be processed.

A sketch of the implemented QMF filter bank for the eight-band SBC (transmitter), including the subband processing sequence according to PATH, is shown in Fig. 12. For example, during cycle 3 (PATH=3) a sample of subband 5 (2000-2500 Hz) is created by respectively passing through HIGH1 in the first QMF stage, HIGH3 in the second stage and LOW7 in

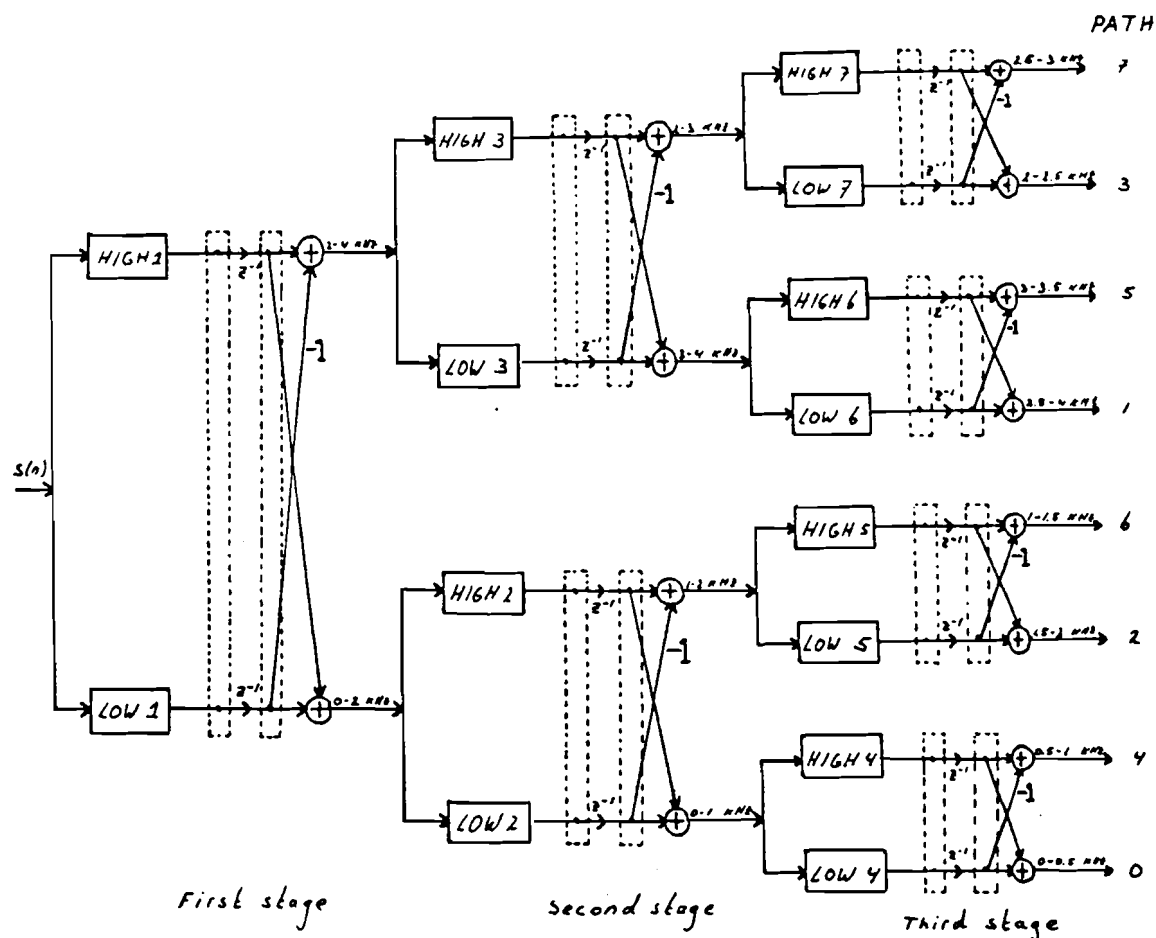


Fig. 12: QMF filter bank implementation for the eight-band SBC (transmitter)

the third stage. Also it has to be taken into account that after high-pass filtering the input signal into the 2-4 kHz part and sampling rate reduction, the 2-3 kHz part is obtained by high-pass filtering in stead of low-pass filtering. This can be verified easily by an examination of the spectra. Similar circumstances occur at some other parts in the QMF tree. As can be seen from Fig. 12, for PATH=0 to 7 the subband processing sequence is : 0-500 Hz, 3500-4000 Hz, 1500-2000 Hz, 2000-2500 Hz, 500-1000 Hz, 3000-3500 Hz, 1000-1500 Hz, 2500-3000 Hz, being respectively subband 1,8,4,5,2, 7,3,6.

For the first QMF stage 32-tap FIR filtering is used (16 taps for the upper branch and 16 taps for the lower branch). For the second and third stage 16- resp. 12-tap FIR filtering is used. The coefficients, as obtained by the Hooke and Jeeves optimization method [21] are depicted in Table 2.

Table 2: Filter coefficients for 12-, 16- and 32-tap QMF filters.

<u>12-tap</u>	
h(0) =h(11)=-0.3809699E-2	h(3) =h(8) =-0.8469594E-1
h(1) =h(10)= 0.1885659E-1	h(4) =h(7) = 0.8846992E-1
h(2) =h(9) =-0.2710326E-2	h(5) =h(6) = 0.4843894E+0
<u>16-tap</u>	
h(0) =h(15)= 0.1050167E-2	h(4) =h(11)=-0.9666376E-2
h(1) =h(14)=-0.5054526E-2	h(5) =h(10)=-0.9039223E-1
h(2) =h(13)=-0.2589756E-2	h(6) =h(9) = 0.9779817E-1
h(3) =h(12)= 0.2764140E-1	h(7) =h(8) = 0.4810284E+0
<u>32-tap</u>	
h(0) =h(31)= 0.6910579E-3	h(8) =h(23)=-0.4187483E-2
h(1) =h(30)=-0.1403793E-2	h(9) =h(22)=-0.3123862E-1
h(2) =h(29)=-0.1268303E-2	h(10)=h(21)= 0.1456844E-1
h(3) =h(28)= 0.4234195E-2	h(11)=h(20)= 0.5294745E-1
h(4) =h(27)= 0.1414246E-2	h(12)=h(19)=-0.3934878E-1
h(5) =h(26)=-0.9458318E-2	h(13)=h(18)=-0.9980243E-1
h(6) =h(25)=-0.1303859E-3	h(14)=h(17)= 0.1285579E+0
h(7) =h(24)= 0.1798145E-1	h(15)=h(16)= 0.4664053E+0

These coefficients are best represented in Q16 format (see 2.3.4.). In the beginning all time varying signals in the QMF filter bank were decided to be represented in the same format as the input signal $s(n)$, i.e. in Q15 format. Later, as will be discussed in 3.5., these representations have been adapted according to the properties of speech spectra. For reasons mentioned in 2.3.2., all filter coefficients (only 30, thanks to the symmetry property, eq. (2), of the QMF's) are permanently available in data memory. The same holds for the delay variables of the first QMF stage and furthermore for all buffer variables. However, the delay variables of the second and third QMF stage are stored in program memory. See also [7, pp. 39-40]. Appendix C contains a listing of the source file of the transmitter QMF bank, appendix H that of the receiver QMF bank.

3.4. Evaluation of the filter bank performance

Proper working of the analysis and synthesis filter bank implementation has been tested with the configuration as shown in Fig. 13.

The analog input, $s(t)$, is bandlimited and subsequently converted to $s(n)$ by 8 kHz sampling followed by A/D conversion. Next the transmitter DSP, loaded with the QMF analysis

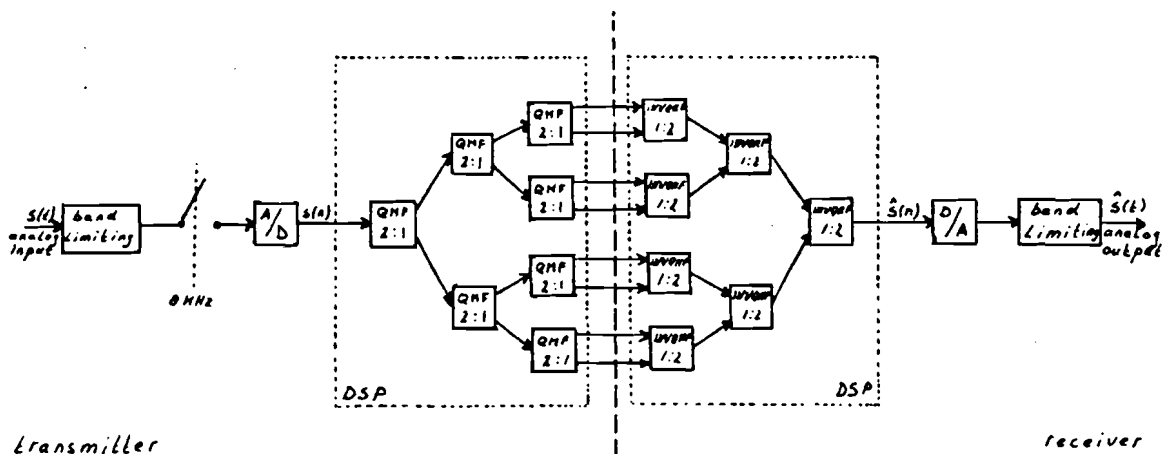


Fig. 13: Test configuration for the analysis and synthesis filter bank.

program, creates for each 8 kHz $s(n)$ sample a 1 kHz subband sample. Each subband sample is transported in a parallel way from the transmitter DSP to the receiver DSP. In the receiver DSP, loaded with the QMF synthesis program, for each incoming subband sample a 8 kHz output sample, $\hat{s}(n)$, is produced. By D/A conversion followed by bandlimiting, $\hat{s}(n)$ is converted to the analog output $\hat{s}(t)$.

How the bandlimiting and the A/D and D/A conversion are performed, will be discussed in chapter 7. The hardware realization and synchronization of the parallel communication between transmitter and receiver system will not be discussed, as it is not functional for the ultimate SBC realization.

The transmitter DSP program is controlled by interrupts. After each 8 kHz interrupt, created by the A/D converter, an interrupt service routine is executed. In this interrupt service routine a sample $s(n)$ is read in and a created subband sample is written out. When the interrupt service routine has finished, the QMF program (for one out of eight paths) is carried out followed by a wait cycle for the next interrupt. QMF program plus wait cycle together form the main routine.

The receiver DSP program is also controlled by 8 kHz interrupts from the regained system clock, which in our circumstances is retained from the transmitter system (see also chapter 7). In the interrupt service routine a subband sample is read in and a reconstructed speech sample, $\hat{s}(n)$, is written out. In the subsequent main routine, the inverse QMF program (for one out of eight paths) is carried out followed by a wait cycle for the next interrupt.

So far the same QMF principle also had been realized during previous work [7,8]. Therefore objective measurements of the filter bank performance give the same satisfactory results (for the reconstruction a "flat" curve in the band of interest) as are extensively discussed and illustrated in [7] and [8].

A subjective test has been carried out by applying an analog

speech signal from a tape recorder to the transmitter input, and listening to the reconstructed speech signal via the audio output of the receiver system. By short-circuit of $s(n)$ and $\hat{s}(n)$, processed speech could be easily compared with unprocessed speech. The intelligibility of processed and unprocessed speech was quite the same. However, processed speech appeared to suffer more from quantization noise, due to the 16-bit representation used in the DSP's.

3.5. Improvement of the subjectively perceived quality

Till now, each time varying signal in the QMF tree has been represented in Q15 format. Watching a long-term speech spectrum as shown in Fig. 14, it is not difficult to see that in general the amplitude range will not be the same for each time varying signal in the QMF tree. So, as mentioned in 2.3.4., the quantization noise can be reduced by choosing an appropriate representation format for each signal in the QMF tree.

Extensive testing, being a very time consuming affair, resulted in the signal representation as depicted in Fig. 15. The signals in the receiver have to be represented in cor-

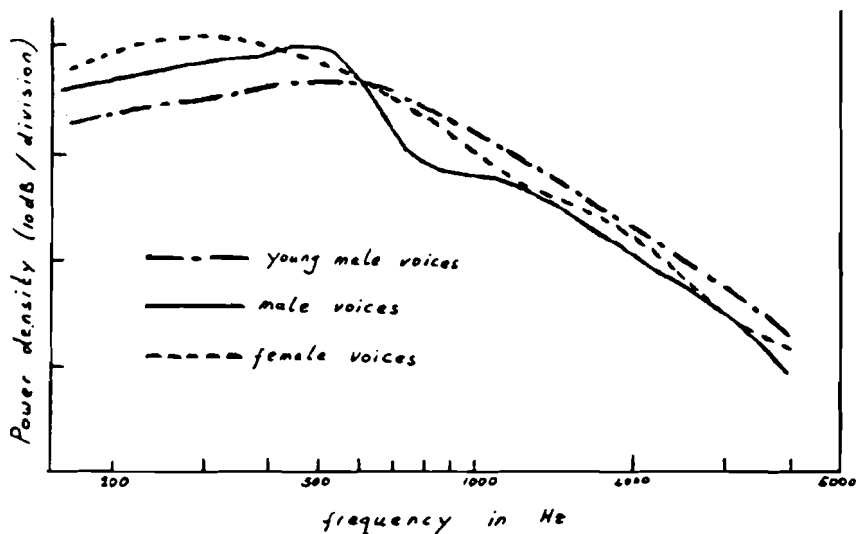


Fig. 14: Long-term spectrum of speech based on measurements by Beranek, Dunn and White [15].

responding formats. This representation has been implemented, performing a reconstructed speech quality that differs not or only marginally from unprocessed speech.

This adjustment of the subband signal representation will also turn out to be advantageous for the encoding of the subband signals (chapter 5), as the dynamic ranges to be covered are better specified.

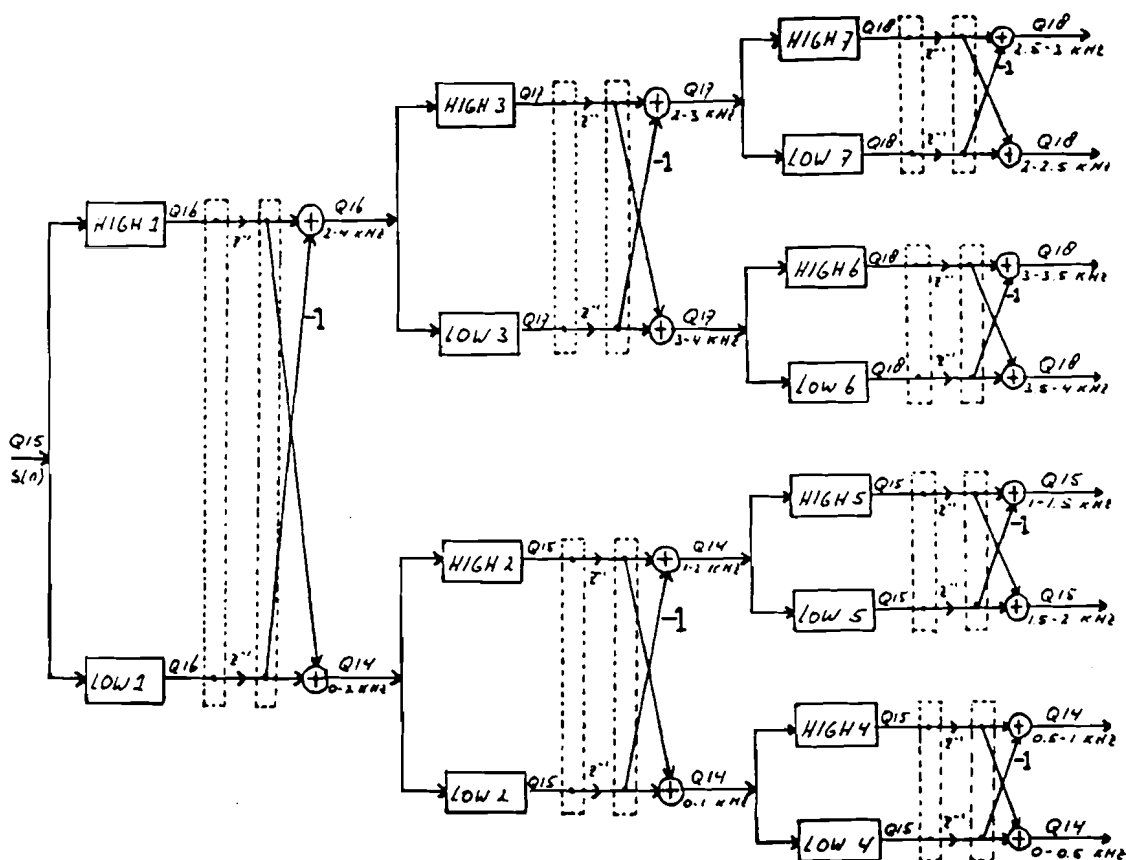


Fig. 15: Signal representation in (transmitter) QMF filter bank.

3.6. Conclusion

The QMF technique has proved to be very suitable for the implementation of the eight-band filter bank. We do not exaggerate unduly if we call $\hat{s}(n)$ a delayed replica of $s(n)$. However, the encoding of the subband signals into a 16 kbit/s bit stream, left out of consideration till now, will cause a performance degradation by adding an amount of quantization noise to the replica of $s(n)$.

4. BIT-ALLOCATION

4.1. Introduction

As discussed in 2.2., after the subband splitting each subband signal is coded independently into a number of bits. The ultimate bit rate has to be 16 kbit/s. So, if we reserve 1 kbit/s for side and synchronization information (see also chapter 6), 15 bits per ms are left to divide over the eight 1 kHz subband signals. This division of 15 bits over the eight subbands is determined by the bit-allocation. For realizing a proper quality it is desirable [6] that the number of bits, allocated to a subband, agrees with its perceptual contribution to the overall subjective quality. As can be seen from Fig. 14, this contribution will not be the same for each speaker. The same holds for one speaker at various moments. In other words, the bit-allocation, which is optimal at one certain moment and for one certain speaker, is not necessary the best at another moment and/or for another speaker.

In this chapter the design and implementation of the bit-allocation algorithm that has been chosen for our purpose is discussed. Also described is an experiment to test the functioning of the implemented algorithm.

4.2. Design considerations

Three methods exist [6] for allocating the bits to the several subbands.

For the first one, called fixed-bit-allocation, the number of bits allocated to a subband is the same at each moment for each speaker. This method will not perform an acceptable quality [3,23].

The second method, called dynamic bit-allocation, is a fully adaptive bit-allocation algorithm, where the power is measured in each band and the bits are successively allocated to the subbands [23]. Proper implementation of this method (in combination with the right coding algorithm) will lead to

very acceptable results. However, this method requires a lot of side information and delay and furthermore is far too complex to fit the DSP device capabilities.

Therefore a compromise solution, called semi-adaptive bit-allocation, has been chosen for the realization of the SBC. A semi-adaptive bit-allocation scheme is, compared to a fully adaptive bit-allocation method, relatively simple to implement, yet gives a significant improvement in performance over a fixed-bit-allocation scheme. Furthermore the required side information is reduced greatly and, as we shall soon see, the semi-adaptive algorithm does not introduce any additional delay over and above that of the QMF transmitter filter bank.

The first QMF filter outputs are used to obtain energy estimates for the bit-allocation algorithm. The spectral envelope estimate, approximated by the short-term-average magnitude, is computed during a window, determined by the average period for which speech signals are stationary. This is done for the speech in each of the subbands, 0-2 kHz (L) and 2-4 kHz (H). At time intervals, dictated by framing considerations (chapter 6), the ratio between average magnitude estimates for the two bands is used to form a three way decision as to whether the speech is voiced, unvoiced or intermediate. The following voicing decision scheme, recommended in the literature [2,3,4], is used:

$L/H > 20$	voiced
$1.5 < L/H < 20$	intermediate
$L/H < 1.5$	unvoiced

Hysteresis is included in the decision-making process to prevent rapid strategy changes due to marginal decisions. If the voiced strategy is already in use, L/H must be less than 10 to change back to the intermediate strategy, and for the unvoiced strategy in use, L/H must be greater than 3 to revert back to the intermediate strategy.

The frequency range 200-3200 Hz is seen as the band of

interest [6,20] in a speech signal. The 0-200 Hz part contains a lot of power that however is not important for the intelligibility of speech. This part can even better be omitted by the bandlimiting filter (Fig. 13) before the QMF splitting takes place, since it has a wrong influence on the voicing decision and causes a lot of aliasing, both resulting in an inefficient coding of the subband signals. Also the part above 3200 Hz, containing only very little power (Fig. 14), is unimportant for speech intelligibility and may be omitted by the bandlimiting filter (as this part contains very little energy, the filtering required at this edge is not as sharp as at the 200 Hz edge). To cover the band of interest as best as possible, the bit-allocation assignment for the three possible voicing strategies as shown in Table 3 is used [2,3,4].

 Table 3: Bit-allocation assignment.

band (kHz)	0.0	0.5	1.0	1.5	2.0	2.5	3.0	3.5
to	0.5	1.0	1.5	2.0	2.5	3.0	3.5	4.0
voiced	4	4	3	2	2	0	0	0
intermediate	4	3	2	2	2	2	0	0
unvoiced	2	3	3	3	2	2	0	0

As can be seen from this table, no bits are allocated to the 3-3.5 kHz subband (i.e. 3000-3200 Hz). Using these bits to improve the coding of the other subband signals, results in a better reconstructed speech quality than inclusion of coding of the 3000-3200 Hz part will do. By consequence it is not necessary to split (and in the receiver to reconstruct) the 3-4 kHz band in the QMF filter bank. However, the 3000-3200 Hz part may not be omitted by the bandlimiting filter, since it contributes to the voicing characteristic of the speech signals.

Ultimately the voicing strategy used also has to be transmitted to the receiver. As there are only three possibilities, the side information required is modest.

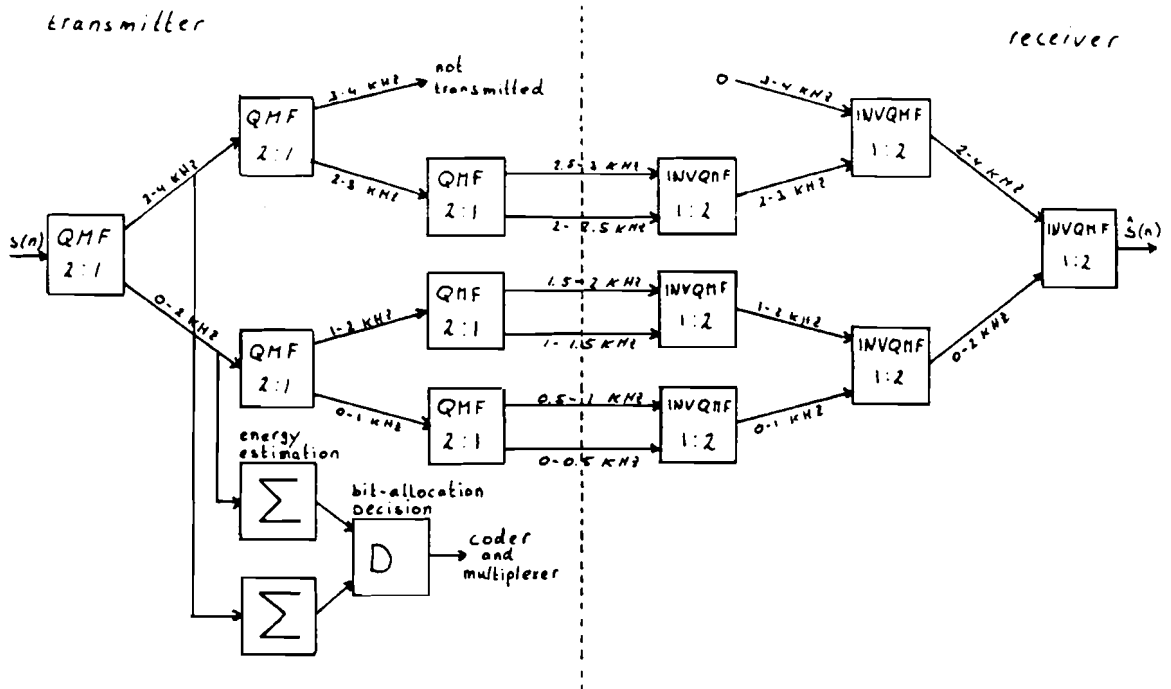


Fig. 16: Semi-adaptive bit-allocation in QMF filter bank.

The filter bank design, extended with the semi-adaptive bit-allocation to be implemented, is sketched in Fig. 16.

4.3. Implementation of the chosen bit-allocation algorithm

The energy summation and the voicing decision itself are straight implementations of the design discussed in 4.2. (using appropriate representation formats for the variables, all being directly stored in data RAM), and therefore need no further explanation. What really matters in the implementation of the bit-allocation algorithm is the synchronization with the rest of the SBC program.

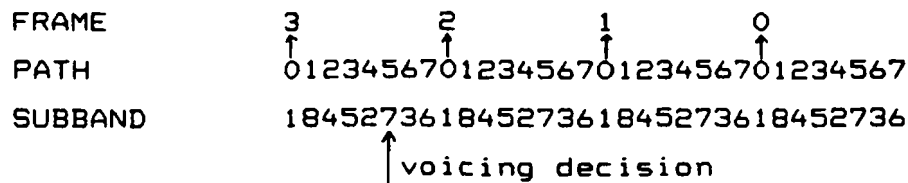
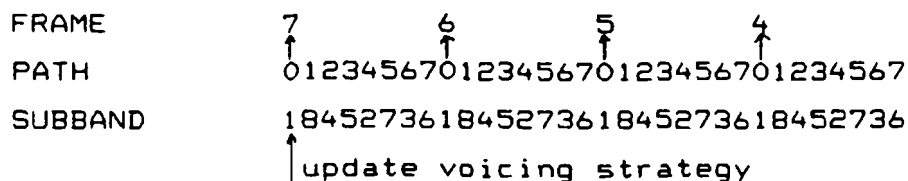
For reasons to be discussed in chapter 6, a frame length of 128 bits (8 ms) is used in multiplexing the coded subband signals. Speech may be considered to be in a stationary state during a period of 8 ms. Therefore, each time a new frame starts, the voicing strategy is updated in agreement with an 8 ms energy measurement. As this energy measurement takes place in the first QMF stage, we have to take into account the processing delay time in the second and third

stage of the QMF filter bank. This delay time is calculated to be 11.5 ms. As the energy measurement takes 8 ms, after the voicing decision it will last another 3.5 ms before the first samples corresponding with the voicing decision enter the PCM coders. Therefore the voicing decision takes place 3.5 ms before a new frame starts, at which the new strategy is taken over (buffering is required). As the 3-4 kHz subband is not transmitted and therefore not further split-
ted in the QMF filter bank, a lot of processing time is saved by executing such voicing decision, where otherwise the 3-4 kHz bandsplitting would occur.

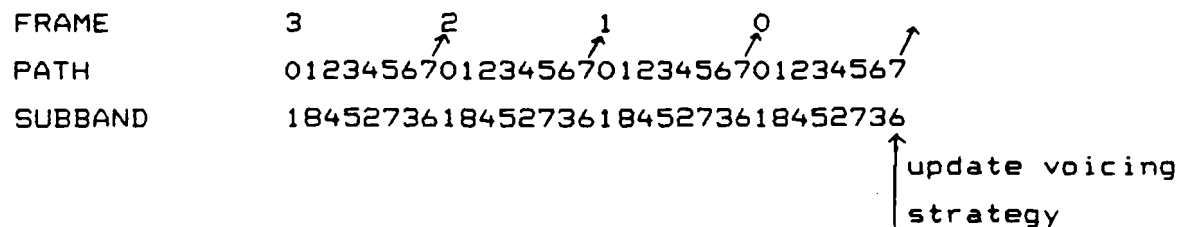
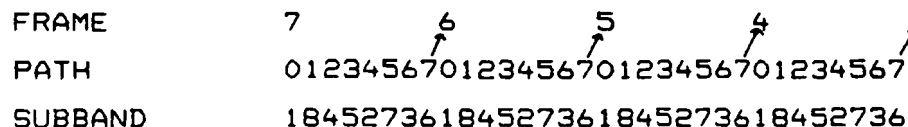
To perform this synchronization a down-counter, called FRAME, is used in addition to the modulo-8 counter PATH, introduced in 3.3. PATH determines the subband to be processed for an 8 kHz cycle. As discussed earlier, for PATH=0 to 7 this is the subband sequence 1,8,4,5,2,7,3,6. For one frame this sequence has to be passed through eight times, kept up by FRAME. PATH is updated for each 8 kHz input, FRAME is updated each time just before a subband 1 sample is coded. At the beginning of a frame, the down-counter FRAME is loaded with the value 7. Each time when PATH=0, FRAME is adjusted. When FRAME=0 and PATH=0 the voicing strategy for the PCM coding is updated and a new frame is started. When FRAME=3 and PATH=5, being ca. 3.5 ms before a new frame start, a new voicing decision is made.

In the receiver the same synchronization principle is used. Here, no voicing decision has to take place. The voicing strategy for the PCM decoding is retained from the side information (available in the frame alignment word, chapter 6). As in the receiver the PCM decoding is done before the QMF reconstruction, FRAME is updated each time that PATH=7. A new frame is started and the voicing strategy is updated when FRAME=0 and PATH=7.

The whole synchronization is depicted in Fig. 17.



(a)



(b)

Fig. 17: (a) Frame synchronization in transmitter
 (b) Frame synchronization in receiver.

The semi-adaptive bit-allocation software has been added to the filter bank programs of appendix C and H.

4.4. Experimental check

After the DSP programs had been extended with the semi-adaptive bit-allocation algorithms, some testing has been done with the same configuration as shown in Fig. 13. Again an analog speech signal was applied to the transmitter input and listening to the reconstructed speech signal took place via the audio output of the receiver. Although the PCM

coding and decoding had not yet been implemented, the perceived quality of the reconstructed speech was not as good as in the case discussed in 3.5. This is a consequence of not transmitting the 3-4 kHz band. However, the quality, for our purpose particularly the intelligibility, was still fairly satisfactory. Therefore, during all further work this quality has been taken to be the maximum quality that can be achieved with our 16 kbit/s speech coder.

This time, also an analog signal, corresponding with the voicing strategy used, was retained from the transmitter system (AIB) and applied to an oscilloscope. In this way the changing of the voicing strategy for different "types" of speech has been confirmed. And indeed changes do not occur within a 8 ms period. Mostly the voicing strategy is constant for more than one 8 ms period, so the choice of 8 ms as a period during which speech signals may be considered stationary is adequate.

Furthermore, by setting breakpoints and single stepping (2.4.) also the synchronization has been checked.

The suitability of the semi-adaptive bit-allocation and the used bit assignments, shown in Table 3, can (apart from 0 bits allocated to the 3-4 kHz band) only be tested if the PCM coding and decoding has been implemented, and therefore will be discussed in 5.4.2.

5. CODING AND DECODING OF THE SUBBAND SAMPLES

5.1. Introduction

Before coding the different subband signals, an incoming sample or difference between an incoming sample and an estimation value of it (in case of differential PCM = DPCM) has to be quantized. Quantization takes place by allocating one out of 2^B quantization levels to the quantizer input, where B is the number of bits assigned for the coding of a subband. Next, the quantized value is coded into B bits. Design of an efficient encoding scheme requires some knowledge of the statistics of the signal to be coded. If we had an a priori knowledge of the statistics of the samples, a nearly optimum quantization scheme would consist of:

- A quantizer matched to the probability density function (PDF) of the samples to be quantized.
- A predictor optimized for the given autocorrelation function of the signal.

In digital speech-encoding systems, we have only a small amount of a priori knowledge of the statistics which, in addition, usually change with time.

- The long-period mean level differs from speaker to speaker.
- At a given mean level, the instantaneous level changes because of variations in speech sounds (c.q. split speech sounds).
- The correlations (as far as present) between successive samples change because of variations in speech sounds (c.q. split speech sounds).

To overcome these problems of unknown statistics, adaptive quantization and (in case of DPCM) adaptive prediction schemes can be used.

We will now discuss the PCM coder and decoder design that best suits our purpose. Then the implementation and testing of it is described.

5.2. Design considerations

Also for the coding of the subband signals several methods are possible. Many of them are extensively discussed in [6, 22, 23, 24, 25].

Concerning the quantization, there are two possibilities, fixed and adaptive quantizers. For reasons like unknown mean level and variations of the instantaneous level, a nonadaptive quantizer will not provide an acceptable quality in our situation. Therefore adaptive quantization has been chosen, in which the 2^B possible quantization levels are not fixed in time. For the adaption two schemes are possible. With the first one, forward adaption, the adaption value is calculated from samples of the input signal. As this has to take place in each subband, this method is far too complex and furthermore requires serial buffering (extra delay) and a lot of side information for transmitting the adaption values to the receiver. The second scheme, backward adaption, is more suitable. Here, the adaption value is calculated from quantized samples. In the receiver the same can be done and therefore no side information is required. Furthermore no extra delay above the filtering delay time occurs, since the buffering required is done in parallel. Hence for our SBC design a backward adaptive quantization scheme (AQB) has been chosen.

For the application of DPCM the following has to be taken into account. With DPCM the complexity of the coder algorithms, especially when the prediction is made adaptive, increases. Furthermore, as there is little or no correlation between the subband outputs [2], the predictor coefficients will be close to zero [26]. So DPCM will perform little or no quality improvement, while the complexity increases. For these reasons it has been decided to use "simple" PCM coding without prediction.

We will now discuss the backward adaptive PCM method, chosen for the (de-)coding of the subband signals, in more detail. See also [6, 7, 15, 26].

A general scheme of a PCMAQB coder is depicted in Fig. 18.

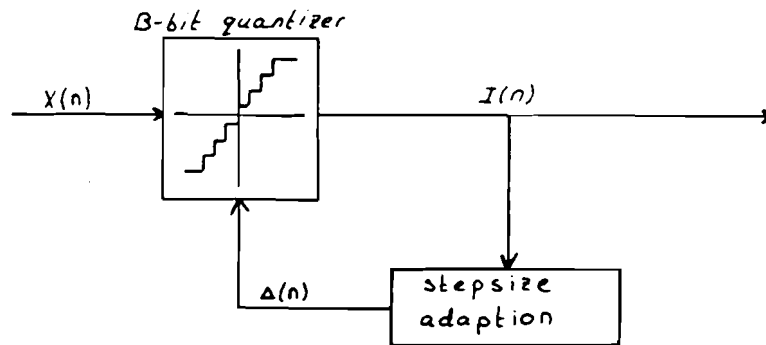


Fig. 18: Block diagram of a PCMAQB coder.

For the quantization of the input signal $x(n)$ a quantizer with a midriser characteristic, as shown in Fig. 19, is used. As can be seen from this figure, a value of $x(n)$ in the range $[I(n) \cdot \Delta(n); (I(n)+1) \cdot \Delta(n)]$ results in a quantized value $\hat{x}(n) = (I(n)+0.5) \cdot \Delta(n)$, where:

$\Delta(n)$ = stepsize (i.e. spacing between quantized levels) at moment n .

$I(n)$ = two's complement value representing one out of 2^B quantization levels, i.e. the B-bit code word, at moment n .

The output of the quantizer is the $\hat{x}(n)$ representing code word $I(n)$. The stepsize adaption strategy used is based on the one-word stepsize memory approach proposed by Jayant, Flanagan and Cumiskey [25]. The robust stepsize adaption is based on the relation:

$$\Delta(n) = [\Delta(n-1)]^\gamma * M(|I(n-1)|) \quad (5)$$

where $\Delta(n)$ is the stepsize used for the encoding at the n th time sample and $\Delta(n-1)$ is the stepsize that was used for the $(n-1)$ th time sample. The value of $\Delta(n-1)$ is raised to a power γ , where $\gamma < 1$ is a coder parameter (to be discussed in more detail in 5.3.). It is then multiplied by a (positive) scale factor $M(\cdot)$ which is a function of the previous code word $I(n-1)$ to give the stepsize estimate $\Delta(n)$. In general,

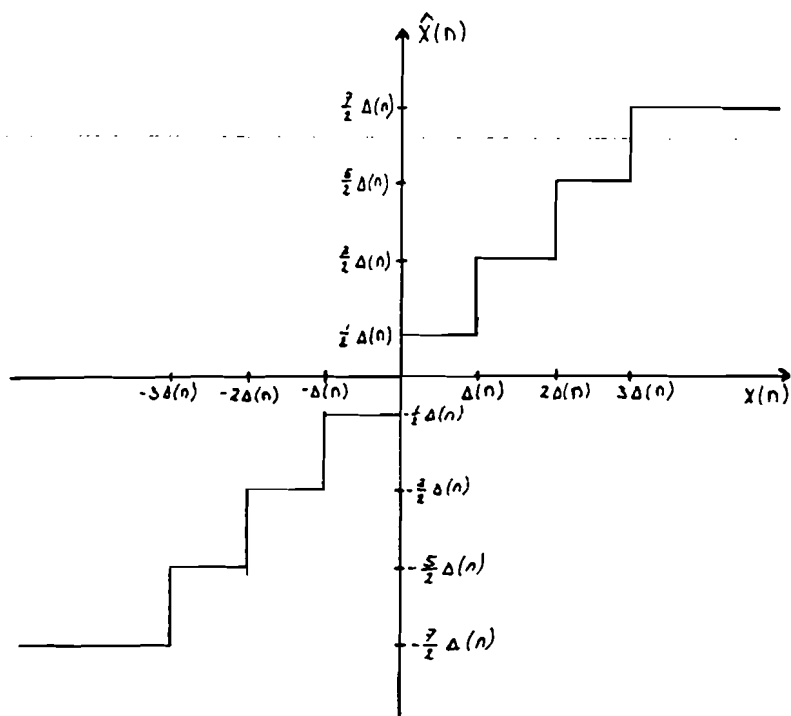


Fig. 19: Uniform quantizer with 8 levels (B=3).

if the previous code word $I(n-1)$ indicates that an upper (absolute) quantizer level was used in encoding, a value of $M(.) > 1$ is used to increase the size of the new stepsize $\Delta(n)$. If $I(n-1)$ indicates that a lower (absolute) amplitude level was used by the quantizer, a value of $M(.) < 1$ is used to reduce the estimation of the new stepsize $\Delta(n)$. Thus the stepsize adaption algorithm is constantly attempting to adjust the stepsize $\Delta(n)$ such that it tracks the rms level of the signal and scales the quantizer characteristic to span the amplitude range of the signal. For practical reasons the stepsize must be limited to the range:

$$\Delta_{\min} \leq \Delta(n) \leq \Delta_{\max} \tag{6}$$

to prevent $\Delta(n)$ to grow beyond the limits of the number representation adopted. The ratio $\Delta_{\max}/\Delta_{\min}$ determines the dynamic range that the coder can handle. In our case a ratio of 2048, ≈ 66 dB, is taken, being within the range of the digital arithmetic. The actual values of Δ_{\min} and Δ_{\max} must be different for each subband (each subband signal is coded

with its own PCMAQB coder), to match properly the dynamic range characteristics of the coders to that of the subband signals, as shown in Fig. 15.

The proportion of the amplitude range that is spanned by the quantizer at a particular time (i.e. for a particular step-size $\Delta(n)$) determines its "loading". If the range of the quantizer is too small relative to the signal range, the quantizer will overload and clip the signal. If it is too large, the quantizer stepsize will be too large, and this will result in an excessive quantization error or noise (often referred to as granular noise). Thus the proper loading of the quantizer is an important factor in maintaining a good reproduction of the signal. The loading is controlled by the choice of the parameters γ and $M(\cdot)$.

In the receiver the same stepsize adaption algorithm is used as in the transmitter, since this adaption takes place according to the coded subband samples. A block diagram of a PCMAQB decoder is depicted in Fig. 20. By consequence of the midriser characteristic of the dequantizer an input $I(n)$ results in the dequantized value $\hat{x}(n) = (I(n) + 0.5) * \Delta(n)$.

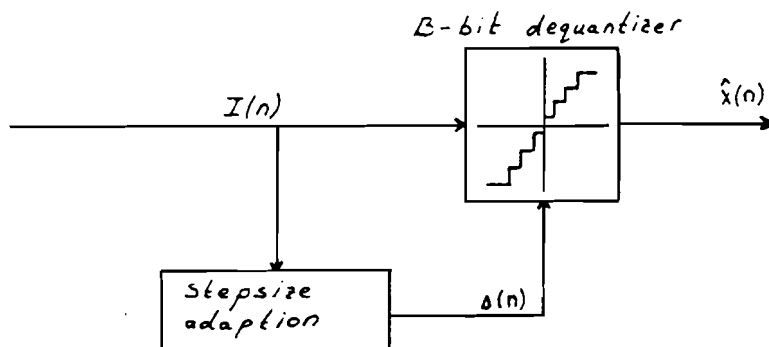


Fig. 20: Block diagram of a PCMAQB decoder.

5.3. Implementation of the chosen (de)coder algorithm

By taking the logarithm, eq. (5) can be written in the form

$$d(n) = \gamma * d(n-1) + m(|I(n-1)|) \quad (7)$$

where:

$$d(n) = \log \Delta(n)$$

$$m(\cdot) = \log M(\cdot)$$

The stepsize adaption is then implemented by the circuit shown in Fig. 21.

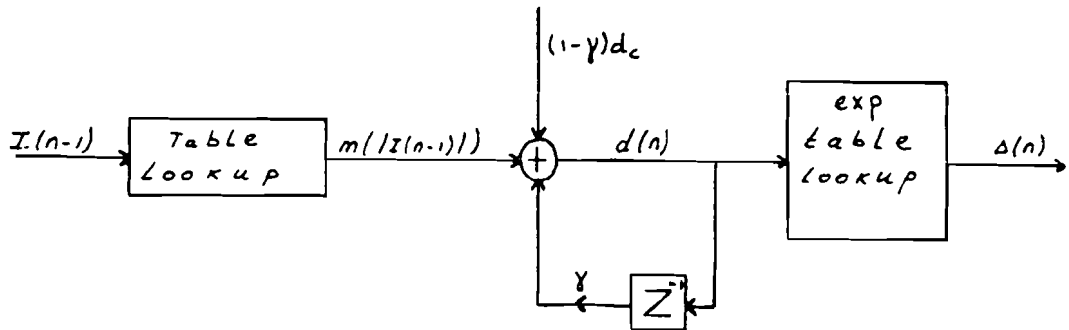


Fig. 21: Stepsize adaption circuit.

The first table lookup converts values of $I(n-1)$ to $m(|I(n-1)|)$, and the second table lookup realizes the exponential conversion from $d(n)$ to $\Delta(n)$. However, as for the quantization the value $1/\Delta(n)$ is needed, in the transmitter an exponential conversion from $d(n)$ to $1/\Delta(n)$ has been implemented. Thus it is seen that the adaption circuit consists of two table lookups and a first-order recursive digital filter which can easily be implemented. The advantage of the method of using table lookups is, that it can be done relatively fast and it can perform functions difficult to calculate in the DSP.

An extra dc input $(1-\gamma)d_c$ is also applied to the circuit in Fig. 21, and it is used to control the loading of the quantizer. Thus eq. (7) has been modified to the form:

$$d(n) = \gamma * d(n-1) + m(|I(n-1)|) + (1-\gamma) * d_c \quad (8)$$

where for d_c has been chosen the practical value of approximately

$$d_c \approx \log(\Delta_{\max}/10) \quad (9)$$

Now it is the right moment to discuss the parameter γ . The adaption leakage factor γ , chosen to be $\gamma=0.99$, forces the realignment of the stepsizes between transmitter and receiver after channel errors occur. Realignment will also occur when the stepsize reaches its maximum or minimum value according to eq. (6), even if γ were chosen to be 1. However, a long time may pass before this maximum or minimum is achieved. Since the cancellation of aliasing in the QMF bank depends strongly on the exact tracking of the stepsizes in each subband, it is therefore preferable to use a value of $\gamma < 1$ to dissipate any effects of channel errors more rapidly. Another effect of the leakage due to γ is that the log of the stepsize, that is $d(n)$ in Fig. 21, tends to decay to zero in the absence of the inputs $m(|I(n-1)|)$ and $(1-\gamma)*d_c$. By adding the term $(1-\gamma)*d_c$ the stepsize toward which $d(n)$ decays can be set to any arbitrary level.

We will now discuss the implementation of the table lookups (stored in program memory) in more detail. In the first table lookup $I(n-1)$ is converted to $m(|I(n-1)|)$ according to:

$$I(n-1) \rightarrow |I(n-1)| \rightarrow M(|I(n-1)|) \rightarrow m(|I(n-1)|) \quad (10)$$

For each subband the same table is used. An appropriate base for taking the logarithms has been found to be 10^4 . Choices for the values of $M(\cdot)$ for different numbers of bits to which the input signal has to be quantized are depicted in Table 4. Experiments have shown that small deviations from the presented values have very little influence on the performance of the adaptive coder. The optimal representation format for the $m(\cdot)$ values is the Q18 format.

The exponential table lookup is not the same for each subband. Each subband has its own dynamic range spanned up by Δ_{\min} and Δ_{\max} . The ratio $\Delta_{\max}/\Delta_{\min}$ is the same (2048) for each subband. The values chosen for Δ_{\max} and Δ_{\min} , which have been adapted to the dynamic ranges of the subband signals, are depicted in Table 5. The six different tables have been

Table 4: PCMAQB coder parameters.

B=	4	3	2
M_1	0.9	0.85	0.85
M_2	0.9	1.0	1.9
M_3	0.9	1.0	
M_4	0.9	1.5	
M_5	1.2		
M_6	1.6		
M_7	2.0		
M_8	2.4		

obtained by taking the logarithm of Δ_{\max} and Δ_{\min} and dividing the range between d_{\min} and d_{\max} into 127 uniformly spaced parts, resulting in 128 entries (convenient for our purpose). For each entry the exponential conversion to Δ (in the transmitter to $1/\Delta$) has been stored in the table. The tables, created in this way, can be implemented very efficiently. If each table is represented in its own optimal format, only one table with 128 entries is needed to represent all six tables.

Table 5: Δ_{\min} and Δ_{\max} for different subbands.

subband	f(Hz)	Δ_{\min}	Δ_{\max}	representation format for Δ
1	0-500	4.882813E-4	1	Q15
2	500-1000	4.882813E-4	1	Q15
3	1000-1500	2.441406E-4	0.5	Q16
4	1500-2000	2.441406E-4	0.5	Q16
5	2000-2500	3.051758E-5	0.0625	Q19
6	2500-3000	3.051758E-5	0.0625	Q19
7	3000-3500	-	-	-
8	3500-4000	-	-	-

For example, Δ_{\min} and Δ_{\max} for subband 1 are optimally represented in Q15 format, resulting in the integer values to be implemented of 16 resp. 32767. For subband 5 Δ_{\min} and Δ_{\max} are optimally represented in Q19 format, which also results in the integer values of 16 resp. 32767. The same holds for intermediate values of Δ . So for each subband this table has to be interpreted with its own representation format and its own input range d_{\min} to d_{\max} . The entries of this table, $d(n)$, are for each subband optimally represented in Q14 format. So this way of implementing allows us to work with six exponential table lookups, matched to the dynamic ranges of the different subband signals, for the price of only one.

Since the six PCMAQB coders, as mentioned above, have a lot in common, they have been all brought together in one PCMAQB coder subroutine, that in some points is passed through differently for the different subband signals. To accomplish this, tables in program memory are used for storing constants as $(1-\gamma)*d_c$, d_{\min} , d_{\max} and the bit-assignments, and for updating variables as $d(n)$ and $I(n)$. These constants and variables are different for each subband and therefore each table contains one entry for each subband. To synchronize with the rest of the SBC program, again the variable PATH is used to determine the subband to be coded. The same holds for the PCMAQB decoders in the receiver.

The essential points in the coding and decoding subroutines, see appendix D resp. G, are discussed below. For a listing of the table lookups is referred to the program memory initializations in appendix B and F.

After the QMF routine, in the transmitter, has delivered a subband sample to the PCMAQB coding routine, the following takes place. First the subband sample is identified using PATH (=0 to 7). This means reading in the number of code bits, dictated by the voicing strategy in use, and reading in the dynamic range (i.e. d_{\min} and d_{\max}) for $d(n)$. Then the stepsize adaption algorithm is executed to determine the new

value of $1/\Delta(n)$ to quantize with. This is done by reading in the previous coded result for that subband, $I(n-1)$, followed by the determination of $m(|I(n-1)|)$, $(1-\gamma)*d_c$ and $\gamma*d(n-1)$. Adding the former three values according to eq. (8) will lead to the new value of $d(n)$. This new value of $d(n)$ is checked on lying within the dynamic range spanned by d_{min} and d_{max} ; if not, saturation to d_{min} or d_{max} takes place. The remaining value of $d(n)$ is saved in its table in program memory. Then $d(n)$ is rounded to one of the entry points of the exponential table lookup, ultimately leading to $1/\Delta(n)$. Having this value of $1/\Delta(n)$ the quantization takes place by multiplying it with the subband sample delivered by the QMF routine and taking the integer part of the product. According to the midriser characteristic of the quantizer this results in the two's complement code word $I(n)$. After checking the range of it, dictated by the number of code bits, and possible saturation the code word $I(n)$ is saved in its table in program memory. Later the multiplexer, to be discussed in chapter 6, will use this table of code words $I(n)$ to create a 16 kbit/s data stream.

In the receiver, the demultiplexer will recover the table of code words $I(n)$ from the 16 kbit/s data stream. This table will be used by the PCMAQB decoder to reconstruct the uncoded subband signals. When this decoding routine is called, first the stepsize $\Delta(n)$ is determined, in the same way as described for the transmitter. Then the subband sample to be decoded, according to PATH, is read in. After adding 0.5 to this value of $I(n)$, multiplication with $\Delta(n)$ takes place, resulting in a dequantized subband sample according to the midriser characteristic. After checking the range of this reconstructed subband sample and possible saturation, it is delivered to the inverse QMF bank for further processing. For more details about the coding and decoding algorithm is referred to the source files, listed in appendix D and G.

5.4. Tests_and_results

For testing the implemented PCMAQB coders and decoders, again the test configuration as shown in Fig. 13 has been used. Only the DSP programs used in this stage differ from those used in Fig. 13. The main routine in the transmitter DSP has been extended with the PCMAQB coding module. After the QMF program has finished, the PCMAQB program is passed-through. Then the wait cycle follows. The same holds for the main routine in the receiver DSP. However, here the PCMAQB decoding is passed-through before the inverse QMF program is called. The configuration of the DSP programs, also including the extensions for the semi-adaptive bit-allocation, is sketched in Fig. 22.

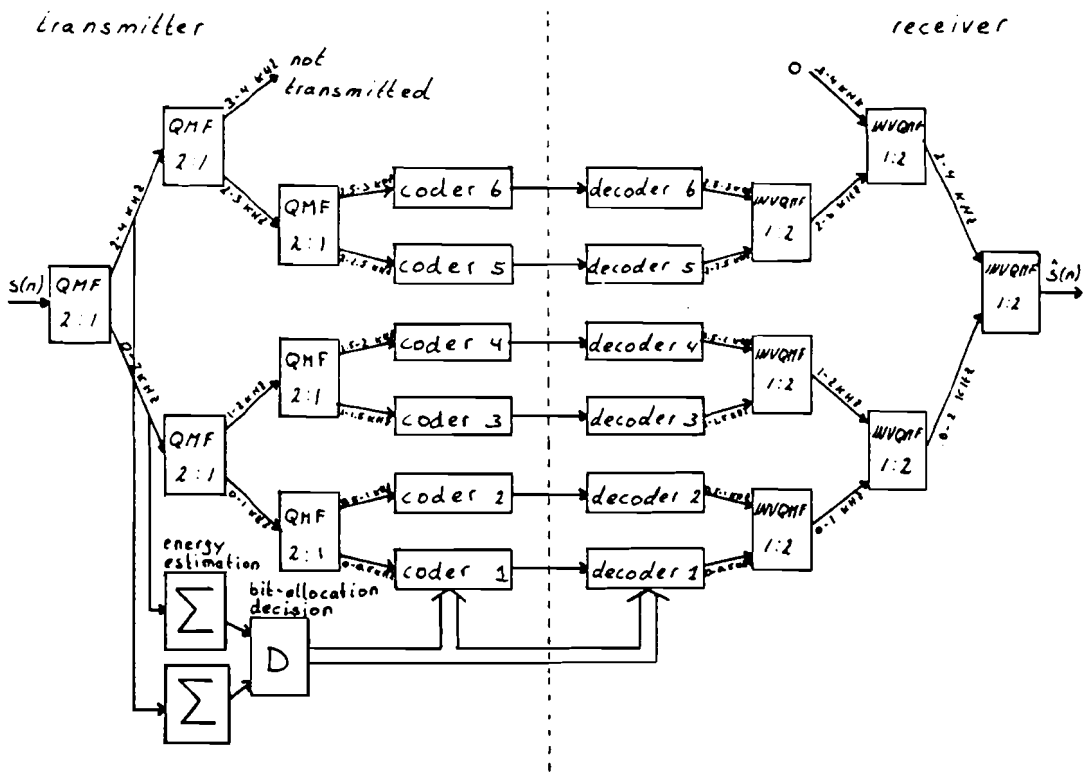


Fig. 22: DSP configuration for testing the PCMAQB coders and decoders.

Again tests with speech signals have been carried out, which is necessary due to the adaptive schemes in the SBC algorithm. In all tests the analog speech input was applied to the transmitter system input, and listening to the recon-

structed speech signal took place via the audio output of the receiver system. For some kind of objective checking, the reconstructed speech signal was also visualized on a memory oscilloscope. Furthermore, the reconstructed speech quality could easily be compared with the maximum quality that can be achieved (up to 3 kHz bandlimited speech quality, see 4.4.) by skipping the PCMAQB coding and decoding (by means of a hardware switch connected with the DSP's).

5.4.1. Performance of the coders and decoders

To confirm the proper working of the coding and decoding itself, the influence of the bit-allocation assignments was eliminated by coding each subband signal with 4 bits. This was allowed, by the fact that the communication between transmitter and receiver system still was performed in a parallel way and not serial. The performance of this 24 kbit/s SBC was rather good. The reconstructed speech differed only very little when compared with uncoded, up to 3 kHz bandlimited, speech.

The loading of the (de-)quantizers was subjectively checked by switching off (in software) the saturation actions in the coding and decoding, since many saturation actions indicate an improper loading of the (de-)quantizers. This resulted in some quality degradation, but not so much as to conclude that the (de-)quantizers were improperly loaded.

Also, the proper loading of the (de-)quantizers has been confirmed objectively, as the behaviour of the reconstructed speech signal showed no excessive clipping on the memory oscilloscope. Finally, also the correctness of the code words in the tables of the transmitter and receiver programs was verified.

After the proper working of the coder and decoder principle had been confirmed in this way, the coding into 16 kbit/s, i.e. the influence of a bit-allocation different from 4 bits per subband sample, could be tested.

5.4.2. Performance of the chosen bit-allocation method

First the semi-adaptive bit-allocation method with the assignments as shown in Table 3, (eliminated in the former testing) was restored. Then the same tests, as described in 5.4.1., have been carried out. The intelligibility of the reconstructed speech was fairly satisfactory. However, a quality degradation, compared to uncoded up to 3 kHz bandlimited speech, was noticeable. Furthermore, tests with several listeners have been carried out to compare the semi-adaptive bit-allocation with fixed-bit-allocation and also to try out other bit assignments. Notwithstanding the fact, that in some cases differences were difficult to perceive, the semi-adaptive bit-allocation with the assignments from Table 3 appeared to provide the best overall performance. However, the difference with fixed-bit-allocation was less than expected from the argument given in 4.2.

Finally the loading of the (de-)quantizers has been investigated, which proved to be proper. This is not only a consequence of a convenient stepsize adaption algorithm but also due to an appropriate representation of the subband signals (see 3.5.) and the use of "six different" exponential table lookups (5.3.). In the development work, preceding the ultimate realization of the coding and decoding described above, also tests have been carried out, where all subband signals were represented in the same Q15 format and were coded using the same exponential table lookup for each subband signal. This resulted in a reconstructed speech signal that excessively suffered from audible clipping and "clicks" due to an improper loading of the (de-)quantizers. Also a dynamic range for the exponential table lookups smaller than 66 dB has been tried out, resulting again in an improper loading of the (de-)quantizers. By the way, due to this dynamic range, the different values of Δ_{min} are such, that in periods of silence ($I(n)=0$ and $\Delta(n)=\Delta_{min}$; reconstructed subband sample $= (0+0.5)*\Delta_{min}$) $1/2 \Delta_{min}$ causes no audible output upon reconstruction.

5.5. Conclusion

The 16 kbit/s coding and decoding of the subband samples, using 2,3 or 4 bit PCMAQB and a semi-adaptive bit-allocation algorithm, results in adding an amount of quantization noise to the replica of $s(n)$ (3.6.). However, the perceived quality of the reconstructed speech signal, especially the intelligibility, is considered acceptable for our applications. At this stage the actual Subband Coding and Decoding has been realized. The only thing left to be described is the realization of the multiplexing and demultiplexing to perform the serial 16 kbit/s communication.

6. MULTIPLEXING AND DEMULTIPLEXING OF THE CODED SUBBAND SAMPLES

6.1. Introduction

To create, after quantization and coding, the ultimate serial 16 kbit/s data stream, the subband signals together with the side information and synchronization bits are multiplexed. This multiplexing has to be performed in a controlled way; by demultiplexing it has to be possible to recover the coded subband samples, the side information and the synchronization bits. This can be accomplished by multiplexing the data to be transmitted into a repetitive framed sequence. Important features, coupled with the multiplexing and demultiplexing are:

-Synchronization between transmitter and receiver, i.e. frame alignment.

i) The time for alignment with 99% probability must not seriously disrupt the speech communication. It is usually a compromise between the time taken to confirm the presence of the framing pattern and the risk of incorrectly aligning to random imitations of it.

ii) False indications of lost frame alignment for an error rate of ca. $1:10^3$ must not occur too frequently.

-Frame organization.

The composition of a frame in the 16 kbit/s serial data stream has to match the framing, i.e. the subband processing sequences and the voicing strategy changes (Fig. 17) in the "actual" (without (de)multiplexing) Subband Coder and Decoder algorithms. This to avoid excessive buffering.

The bit synchronization in the demultiplexer is assumed to be located in a (higher order) part outside the Subband Decoder algorithm, where the demultiplexing of the Subband Coder signal (16 kbit/s) and another signal (see 1.) has to be done. Therefore the bit synchronization will not be described here.

In this chapter the design, implementation and testing of the multiplexer and demultiplexer for our SBC will be discussed.

6.2. Design considerations

Based upon recommendations from literature [3,27] and previous work [6,8,28] it has been decided to use the following strategies for the framing and synchronization.

To acquire frame alignment, when the first bits are received or alignment has been lost, it is for the demultiplexer to search for and recognize the Frame Alignment Word (FAW), present in a fixed position in each frame. Then the demultiplexer has to lock its timing counters into the correct phase relationship with the incoming signal, and examine the FAW in two successive confirmatory frames. Non-recognition of the FAW in its expected position in either of these two frames causes a recommencement of the search. This is done to safeguard against false alignment by imitations of the FAW within the signal.

The minimum time that is necessary to acquire and confirm frame alignment is, therefore, between 2 and 3 frame periods, depending on the point within the frame that a valid signal is applied and the search commences. However, the incoming signal contains an essentially random occurrence of ones and zeros in the information digit time-slots, and there is a probability of these digits imitating the FAW. The probability of an imitation in any position is 2^{-m} for a random bit stream, m being the number of bits in the FAW. Such imitations of the FAW will lead to an increase of the time required to acquire frame alignment, due to the greater number of false starts; the false alignment being rejected at the first or second confirmatory frame.

The time, t_1 , to acquire frame alignment with a 99% probability of not being exceeded can be estimated using the following simplified formula for the alignment strategy described above:

$$t_1 = F(3 + (F-1)2^{-m} + 2*3[(F-1)2^{-m}]^{4^F}) / 16 \quad \text{milliseconds} \quad (11)$$

where:

F is the total number of bits in a frame of the multiplexed signal.

m is the number of bits in the FAW.

To avoid false alarms due to bit errors as much as possible, frame alignment is considered to have been lost when 4 consecutive FAW's are incorrectly received in their predicted positions. The probability of random digital errors causing this condition to be fulfilled is approximately $(mp)^4$ when p, the bit error rate, is low; better than, say, $1:10^3$. It follows that the mean time between false losses of frame alignment for a given error rate is,

$$t_2 = F / (16(mp)^4) \quad \text{milliseconds} \quad (12)$$

To meet the requirements i) and ii) in 6.1., t_1 has to be small (order of ms) and t_2 has to be large (order of hours). Appropriate values for the frame length F and the FAW length m have been chosen to be 128 resp. 8 bits. In that case the frame duration is 8 ms and the information bit rate is 15 bits per ms. These choices have already been used in designing the bit-allocation and synchronization for the "actual" Subband Coding and Decoding (chapter 4).

With respect to the organization of a frame, the following can be said. As mentioned above, for synchronization reasons the first 8 bits of such an 128 bits frame contain a FAW. As the side information, required for our SBC design, only concerns the bit-allocation (3 possibilities) used for the coding of the subband samples during a whole frame period, this is inserted in the FAW. So three FAW's are possible, each being composed in a special way to reduce the probability of imitating it:

FAW1 = >76	voiced
FAW2 = >96	intermediate
FAW3 = >CC	unvoiced

The remaining 120 bits are all used for transmission of the coded subband samples. These 120 bits are filled in, taking into account the way of subband processing in the "actual" SBC, see Fig. 17.

6.3. Implementation of the chosen multiplexer and demultiplexer algorithm

Before describing the implementation of the multiplexer and demultiplexer themselves, we will discuss the application of the 16 kHz clock signal, needed to create a 16 kbit/s bit stream. As discussed earlier, the "actual" SBC algorithms are controlled by an 8 kHz interrupt signal, updating the timing variables PATH and FRAME. Therefore, to complete the overall SBC system, the DSP programs have been adapted as follows.

The transmitter DSP is controlled by interrupts, from now on a 16 kHz clock signal. Each 16 kHz interrupt created by the A/D converter (see chapter 7) causes an execution of the interrupt service routine. In this interrupt service routine a sample of $s(n)$ is read in, followed by a call of the multiplex subroutine to create one bit to be written out. The main routine is passed-through for every other 16 kHz interrupt (let say the odd ones), realizing the same 8 kHz control as before. In this main routine, thus once interrupted by a 16 kHz interrupt, the QMF filtering and PCMAQB coding is carried out, followed by a wait cycle for the next odd interrupt.

The receiver DSP program is controlled by 16 kHz interrupts from the regained system clock, for which again the transmitter system clock has been used. In the interrupt service routine the demultiplex subroutine is called to read in one bit from the 16 kbit/s stream and furthermore a reconstructed speech sample, $\hat{s}(n)$, is written out. In the main routi-

ne, to be passed-through for every other interrupt, the PCMAQB decoding and the QMF reconstruction takes place, followed by a wait cycle for the next odd interrupt. However, when the receiver is not in alignment $\hat{s}(n)$ is set to zero, without calling the PCMAQB decoding and QMF reconstruction routines.

Furthermore, considering the 16 kHz transmitter input, $s(n)$, only one out of two inputs is used in the main routine. In the receiver, the reconstructed output from the inverse QMF filtering is a sample sequence with a rate of 8 kHz. To reduce the aperture effect [14, pp. 302-304] this sequence is interpolated to obtain the ultimate output signal, $\hat{s}(n)$, with a sampling rate of 16 kHz.

To accomplish the whole synchronization, beside the variables PATH and FRAME a variable called STATE is used. STATE is a modulo-16 counter, being updated for every 16 kHz interrupt. So, for every odd value of STATE, PATH is updated and the main routine (in transmitter or receiver) is passed-through. As discussed in 4.3., FRAME is updated for every sequence of PATH from 0 to 7, which has to be done eight times for one frame period in the "actual" Subband Coding or Decoding.

Now the multiplexer and demultiplexer implementations are discussed in more detail.

6.3.1. Multiplexer

Apart from timing considerations, the multiplex algorithm is not very complicated. Each time the multiplex routine is called, a variable is checked to determine the sort of information, i.e. from the FAW or a certain subband, to be transmitted. A second variable is checked to determine which bit from that FAW or subband has to be transmitted. The transmission of each bit is done via a send buffer. Each time before the first bit of a FAW or coded subband sample has to be transmitted, the send buffer is loaded with that FAW or subband sample code word. The FAW is retained from

the bit-allocation decision network, while the coded subband samples are retained from the "I(n) table" (see 5.3.) in program memory. The variable that indicates the bit to be transmitted is updated for each bit sent. The variable that indicates the sort of information to be transmitted is updated after the last bit of the send buffer has been sent. To match the way of subband processing in the QMF splitting and PCMAQB coding (Fig. 17) first the bits from subband 1 are sent, followed by sending the bits from resp. subband 4,5,2,3 and 6. This sequence is repeated eight times for one frame. The sending of these eight sequences is preceded by sending the FAW of the frame. Furthermore, to ensure that each subband sample is retained from the "I(n) table" on the right moment (i.e. after it has been created and before the next sample of that subband is created) the multiplexer frame periods are delayed three 16 kHz interrupts (being the minimum possible) with respect to the frame periods in the "actual" transmitter SBC. The subband processing in the QMF splitting and PCMAQB coding for one frame period, and its corresponding multiplexer frame for the three voicing strategies resp. voiced, intermediate and unvoiced, are depicted in Fig. 23. In this figure, the numbers in the upper line represent the subband processed in the QMF splitting and PCMAQB coding between two odd 16 kHz interrupt service routines. The numbers in the next line represent the subband

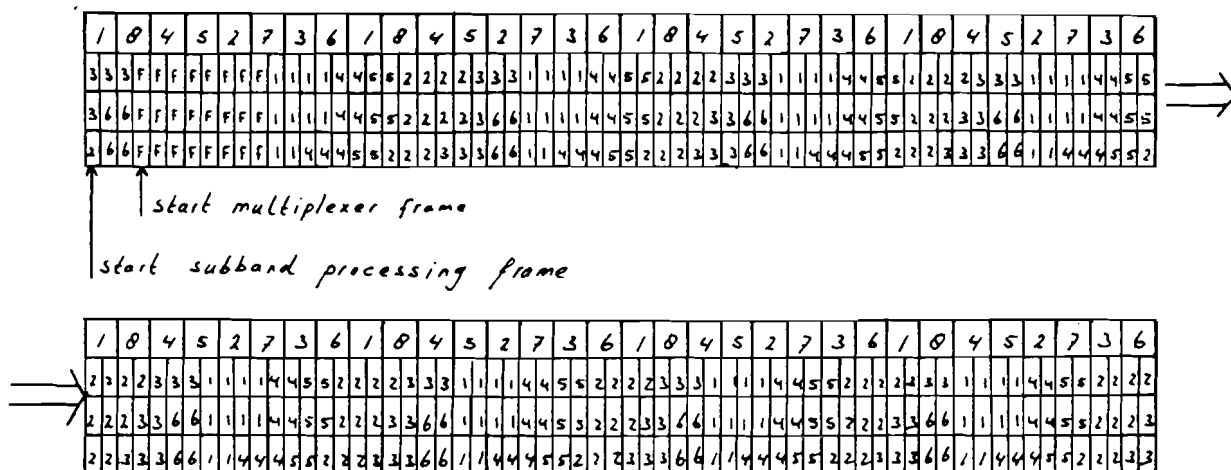


Fig. 23: Frame composition in multiplexer with respect to a subband processing frame.

from which a bit is sent in a 16 kHz interrupt service routine for the voiced bit-allocation strategy; a "F" represents a FAW bit. The same holds for the third and fourth line, but then for the intermediate resp. unvoiced bit-allocation strategy.

The source file of the multiplex subroutine is listed in appendix E.

6.3.2. Demultiplexer

The demultiplex algorithm is more complex, as it also has to take care of the frame alignment. Each time the demultiplex routine is called, it is determined whether the receiver is in alignment or not.

When the receiver is not in alignment, either a search for a FAW takes place or, when a FAW already has been detected, the confirmatory stage is executed. For the FAW search, a received bit is inserted in the receive buffer (a shift register), which is subsequently checked on containing a FAW. If a FAW is detected, then in the next demultiplex call jumping to the confirmatory stage takes place, otherwise again a FAW search is executed. In the confirmatory stage, a received bit is inserted in the receive buffer and, if it is the moment to expect a FAW (128 bits after the previous FAW), the receive buffer is checked on containing a FAW. In the demultiplex calls, jumping to this confirmatory stage repeats until alignment is definitively confirmed or denied. When denied, the next demultiplex call concerns a FAW search. When confirmed, the receiver timing is set by loading the synchronization variables STATE, PATH and FRAME with the appropriate values, and in the next demultiplex call demultiplexing takes place according to the "in alignment situation".

When the receiver is in alignment, the inverse of the multiplex algorithm is executed. Again two variables are used to determine which bit for which subband or FAW is to be received. The received subband samples are stored in the "I(n) table" (see 5.3.) in program memory. The side information

appendix I.

6.4. Experimental confirmation of the proper working of the multiplexer and demultiplexer

With the configuration from Fig. 13, the 8 kHz A/D and D/A conversions replaced by 16 kHz A/D and D/A conversions and the DSP programs updated to perform the functioning depicted in Fig. 25, experiments have been carried out. Again analog speech was applied to the transmitter system input and the audio output of the receiver system was used to listen to the reconstructed speech.

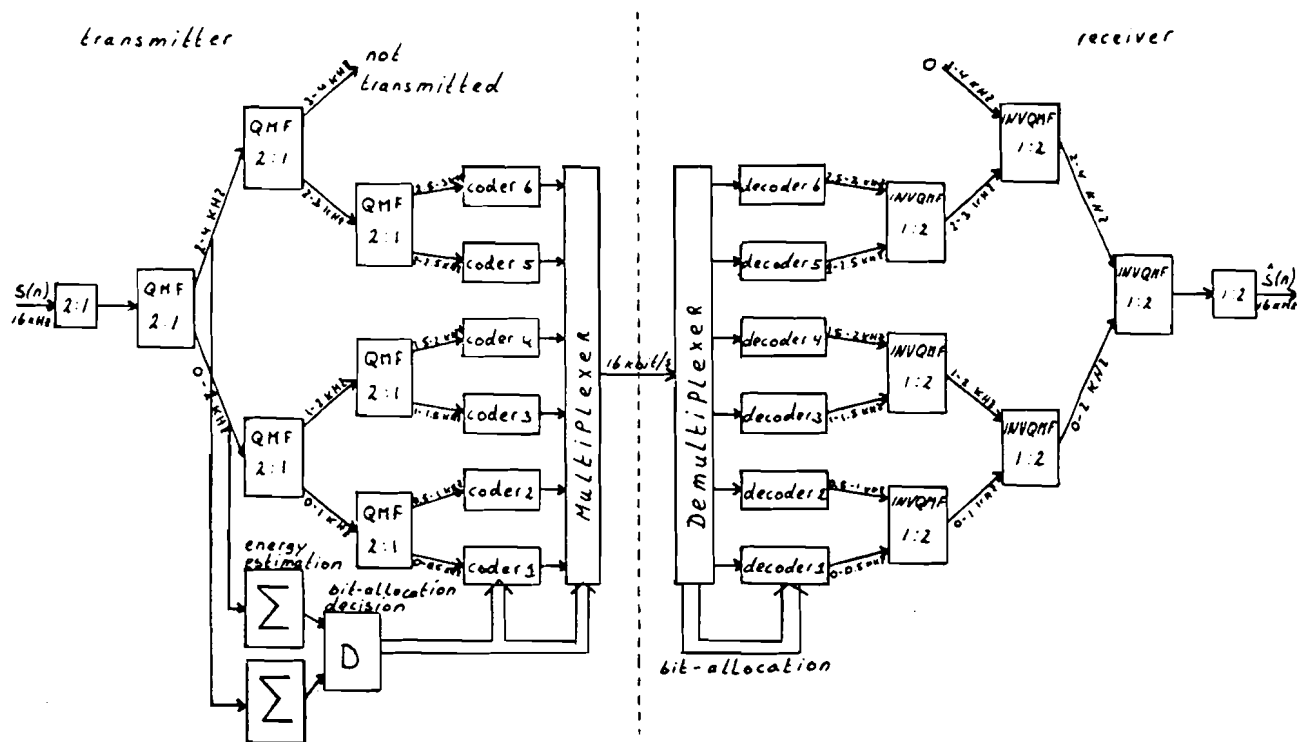


Fig. 25: Block diagram of the implemented Subband Coder and Decoder on two DSP's.

After starting up the transmitter and receiver program, frame alignment in the receiver was acquired very fast (hardly to perceive). When in alignment, the perceived quality obtained with this "serial 16 kbit/s SBC" was exactly the same as that obtained for the "parallel 16 kbit/s SBC" (5.4.2.). Furthermore, proper functioning in situations that cause loss of frame alignment and frame realignment has been

confirmed by imitating these situations (interrupting the communications and/or synchronization).

Also, for these situations, the necessity of an adaption leakage factor $\gamma < 1$ in the PCMAQB (de-)coding, eq. (8), has been proved. Choosing this factor equal to 1 resulted in a lasting mis-adaption between the PCMAQB coding and decoding stepsizes, leading to unintelligible reconstructed speech.

So, also the multiplexer and demultiplexer implementations may be considered to function properly. With that, as these were the last modules to be added, the description of the SBC realization concerning the DSP implementations has finished. Therefore Fig. 25 is the block diagram of the implemented Subband Coder and Decoder.

All the necessary source files are listed in appendix B to E for the transmitter DSP, and in appendix F to I for the receiver DSP.

7. HARDWARE_CONSIDERATIONS

7.1. Introduction

So far the hardware environment of the DSP's, also necessary for a proper working SBC realization, has been left undiscussed.

In this chapter we will discuss the essentials of the band-limiting, the A/D and D/A conversion and furthermore the serial transmission from transmitter to receiver DSP.

For detailed information about the A/D and D/A converters, and for a hardware scheme of the AIB (pin connections etc.) is referred to the AIB book [11].

7.2. I/O_of_the_transmitter_DSP

The 16 kHz 12-bit linear PCM digital input, $s(n)$, to the transmitter DSP, and the 16 kHz interrupt signal that controls the DSP program (6.3.) are created in the following way.

The analog speech signal, $s(t)$, is bandlimited, for which in our case a standard telephone-channel filter (appendix A) has been used. Then the signal is applied to J2 of the AIB, which is via a sample-and-hold-circuit the A/D converter input. Via the software in the DSP program, this A/D converter is programmed to operate with a sampling rate of 16 kHz, and to deliver beside the $s(n)$ signal also a 16 kHz interrupt signal to the DSP. In short the A/D converter works as follows. Upon receipt of a start-of-conversion signal (programmed to be a 16 kHz clock signal) the A/D conversion starts, and after the conversion has been completed, an end-of-conversion signal is created and applied to the interrupt pin of the transmitter DSP. This causes an interrupt (16 kHz) to read in the A/D converted sample and to control the DSP program. The end-of-conversion signal is also used for taking a new sample in the sample-and-hold-circuit. Furthermore, the analog speech signal is supposed to match the dynamic range of the A/D converter, which can be achieved by

means of an amplifier.

The serial 16 kbit/s output stream from the transmitter can be retained from pin 38 of port P1 (a buffered AIB output port). This port P1 is programmed via the DSP program to operate in the sample delay mode, to ensure periodicity in the 16 kbit/s output stream. This means that an output bit from the transmitter DSP is first stored in a primary buffer. By means of a pulse from the same 16 kHz start-of-conversion signal as mentioned above, it is transferred to a secondary buffer, being the interface with the "outside world". As the primary buffer is filled via the DSP program far before a start-of-conversion pulse occurs, it always contains stable data on the moment of transfer.

7.3. I/O of the receiver DSP

In the receiver system, the clock signal that controls the DSP program is not created on the AIB itself as is the case for the transmitter system. In the receiver system an externally regained 16 kHz system clock (to be discussed in 7.4.) controls the synchronization.

The serial 16 kbit/s input stream to the receiver system has to be delivered to pin 12 of port P1. Via the DSP program this AIB input port is programmed to operate in the asynchronous receive mode. By means of a pulse from the regained system clock a bit from the input stream is clocked into an input buffer. As the regained system clock is applied to the interrupt pin of the DSP, this pulse also creates an interrupt to read in the buffered input bit and to control the DSP program.

The 16 kHz digital output samples, $s(n)$, from the receiver DSP are delivered to the D/A converter that is programmed to operate in the transparent mode. Each sample is directly D/A converted without double buffering. Next, the D/A converted signal is bandlimited with a up to 4.7 kHz bandlimiting filter resident on the AIB. Finally, the analog output is available via J3, or, when an audio output is desired, via J4 (AIB outputs).

7.4. Communication between transmitter and receiver system

As mentioned in 6.1., in the receiver the system clock is assumed to be regained in a (higher order) part outside the subband decoder algorithm. Therefore, for our (testing) purposes this clock signal has been retained from the transmitter system (hard-wired clock). The signal taken, is the 16 kHz end-of-conversion signal. An end-of-conversion pulse occurs ca. 25 μ s after a start-of-conversion pulse used in the transmitter system to clock-out a bit of the 16 kbit/s stream. So, when this end-of-conversion signal is used as a regained system clock in the receiver, always stable data are clocked into the input buffer of the receiver system. Furthermore in this way the transmitter and receiver DSP programs are controlled by the same 16 kHz clock.

In the transmitter system, the end-of-conversion signal is retained from pin 9 of U40 on the AIB, and is via a 50 Ω line driver (SN74S140N) applied to a BNC connector. In the receiver system, this end-of-conversion signal is via a BNC connector applied to pin 13 of port P1.

If both pin 38 of port P1 in the transmitter system and pin 12 of port P1 in the receiver system have been connected with a BNC connector too, the communication between transmitter and receiver system is realized by means of two coaxial cables. One for the 16 kHz system clock and one for the 16 kbit/s serial data stream.

With a frequency counter the system clock frequency has been measured to be 15.974 kHz.

7.5. Outline of the hardware configuration that realizes a 16_kbit/s Subband Coder and Decoder

In Fig. 26 a simple diagram of the SBC hardware and its connections is shown. Furthermore, for the AIB's, to perform the functioning as described in this chapter its jumpers have to be set according to the settings depicted in Table 6.

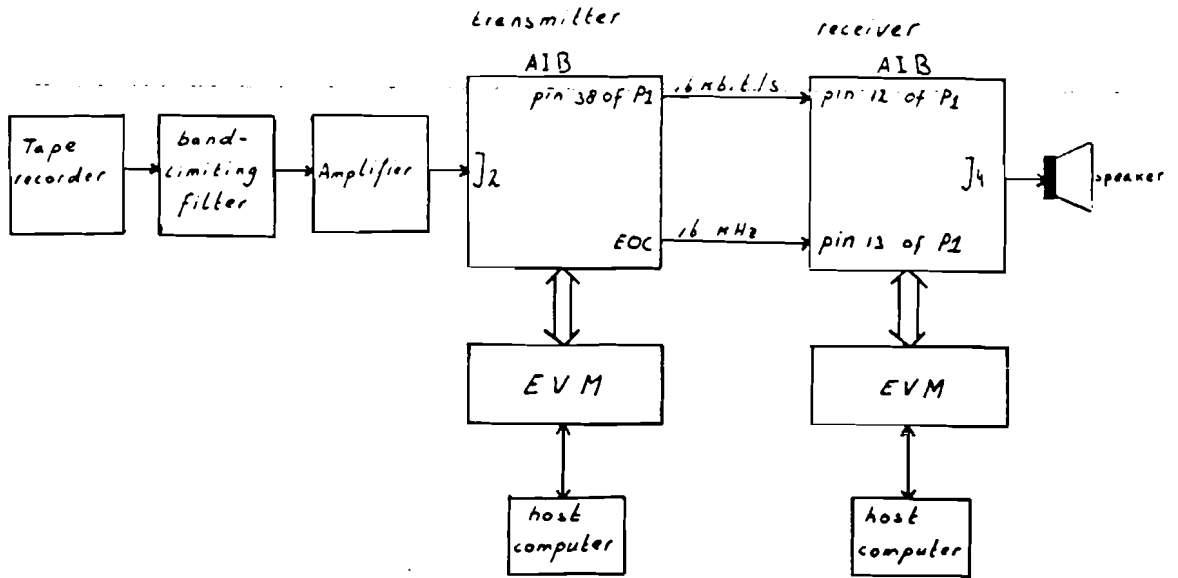


Fig. 26: SBC hardware configuration.

Table 6: AIB jumper settings.

jumper	setting	description
transmitter		
E1	4-5	Connects J2 input jack to A/D (bypasses filter)
E2	don't care	
E3	don't care	
E4	not connected	Leaves Vcc pin on target socket open
E5	2-3	Connects A/D end-of-conversion signal to interrupt pin of TMS32010 emulator socket
E6	1-2	Connects sample and hold to A/D input
receiver		
E1	don't care	
E2	1-2	Connects output filter to J3 output jack
	3-4	Connects D/A converter to output filter
E3	1-2	Connects analog output to audio amplifier
E4	not connected	Leaves Vcc pin on target socket open
E5	2-3	Connects regained system clock to interrupt pin of TMS32010 emulator socket
E6	don't care	

8. UTILIZATION_OF_THE_TRANSMITTER_AND_RECEIVER_DSP

At all stages in the design of the Subband Coder, a prime consideration was the optimum use of the DSP resources. Table 7 illustrates the allocation of processing time for a 0.125 ms period (in which one 8 kHz speech sample is processed) and the utilization of program and data memory. The brackets () indicate the percentage use of the total available resource. Concerning the processing time, the worse case situations are depicted. Table 8 shows the decoder utilization for the inverse processes to restore the original speech signal.

-----Table 7: Transmitter DSP utilization-----

function	processing time instruction cycles/0.125 ms	program memory RAM locations
Initialization	-	391
Main program	19	16
Interrupt service routine	68	34
QMF subroutine	217	668
PCMAQB subroutine	110	110
Multiplex subroutine	80	282
total	494 (79%)	1501 (37%)

Data memory (number of RAM locations)=125 (87%)

----- Table 8: Receiver DSP utilization -----

function	processing time instruction cycles/0.125 ms	program memory RAM locations
Initialization	-	391
Main program	22	23
Interrupt service routine	74	40
Inverse QMF subroutine	230	675
PCMAQB subroutine	117	118
Demultiplex subroutine	105	396
total	548 (88%)	1643 (40%)

Data memory (number of RAM locations)=129 (90%)

Considering processor speed and data RAM utilization, we may conclude that the DSP capacities are almost completely exploited. Using one DSP for the transmitter and one for the receiver, a more complex SBC algorithm can most probably not be realized.

9. SUGGESTIONS FOR (POSSIBLE) IMPROVEMENT OF THE SUBBAND CODER PERFORMANCE

So far, the bandlimiting in the transmitter and receiver occurs via resp. a standard telephone-channel filter and a lowpass filter resident on the AIB. The performance of the SBC will improve when bandlimiting filters, matching the band of interest 200-3200 Hz (4.2.) more properly, are used. These bandlimiting filters can be realized in an analog or digital form. During previous work [7, pp. 45-58] such a digital bandlimiting filter has been developed. However, adding this algorithm to our SBC program will largely exceed the DSP capacity for both transmitter and receiver. And therefore the application of digital bandlimiting filters will require another two DSP's, one for the transmitter and one for the receiver.

One of the reasons for not choosing DPCM for the coding of the subband signals, was the fact that little or no correlation exists between the subband samples. However, according to [29] some correlation exists in the first two subbands (0-500 Hz and 500-1000 Hz). In coding, whitening these two subband signals using DPCM in stead of PCM will probably improve the SBC performance. It is questionable, however, whether the DSP's can cope with the resulting increase of complexity.

Besides the literature referred to so far, also the references [30] to [60] have been studied. Here, some alternative SBC algorithms are presented. However, our SBC implementation still remains the best compromise between quality and complexity for our purposes.

10. CONCLUSIONS

The Subband Coding technique, using QMF filters for the bandsplitting and reconstruction, PCMAQB for the coding and decoding of the subband signals and furthermore applying a semi-adaptive bit-allocation strategy, has proved to be a very appropriate method for the realization of a 16 kbit/s speech coder and decoder on one DSP each.

Reviewing the design constraints, the quality performance of the implemented SBC, especially the intelligibility of the reconstructed speech, is fairly satisfactory. The capacities of the DSP's are almost completely utilized. So, also the TMS32010, selected for implementing the SBC algorithm, has showed to be a suitable choice.

For professional applications (i.e. without development systems) in future, the Subband Coder will consist of the following main parts:

- a digital or analog bandlimiting filter (200-3200 Hz).
- an amplifier to accomplish the matching of the dynamic ranges of the input speech signal and the A/D converter.
- a 16 kHz sample and hold circuit.
- a 16 kHz 12 bit linear PCM A/D converter.
- a TMS32010 DSP; program memory is partly present in on-chip or off-chip PROM (burnt-in with the Subband Coder program) and partly in off-chip RAM (for tables and variables, not fixed in time).
- a power supply.

The main parts of the Subband Decoder will be:

- a TMS32010 DSP; program memory is partly present in on-chip or off-chip PROM (burnt-in with the Subband Decoder program) and partly in off-chip RAM (for tables and variables, not fixed in time).
- a 16 kHz 12 bit linear PCM D/A converter.
- a digital or analog bandlimiting filter (200-3200 Hz).

-an audio amplifier+loudspeaker.

-a power supply.

Since both Subband Coder and Decoder will be part of a satellite communication system (see 1.), their 16 kHz system clocks can be retained from the (64 kHz) master clock available in the transmitter and receiver of the satellite link.

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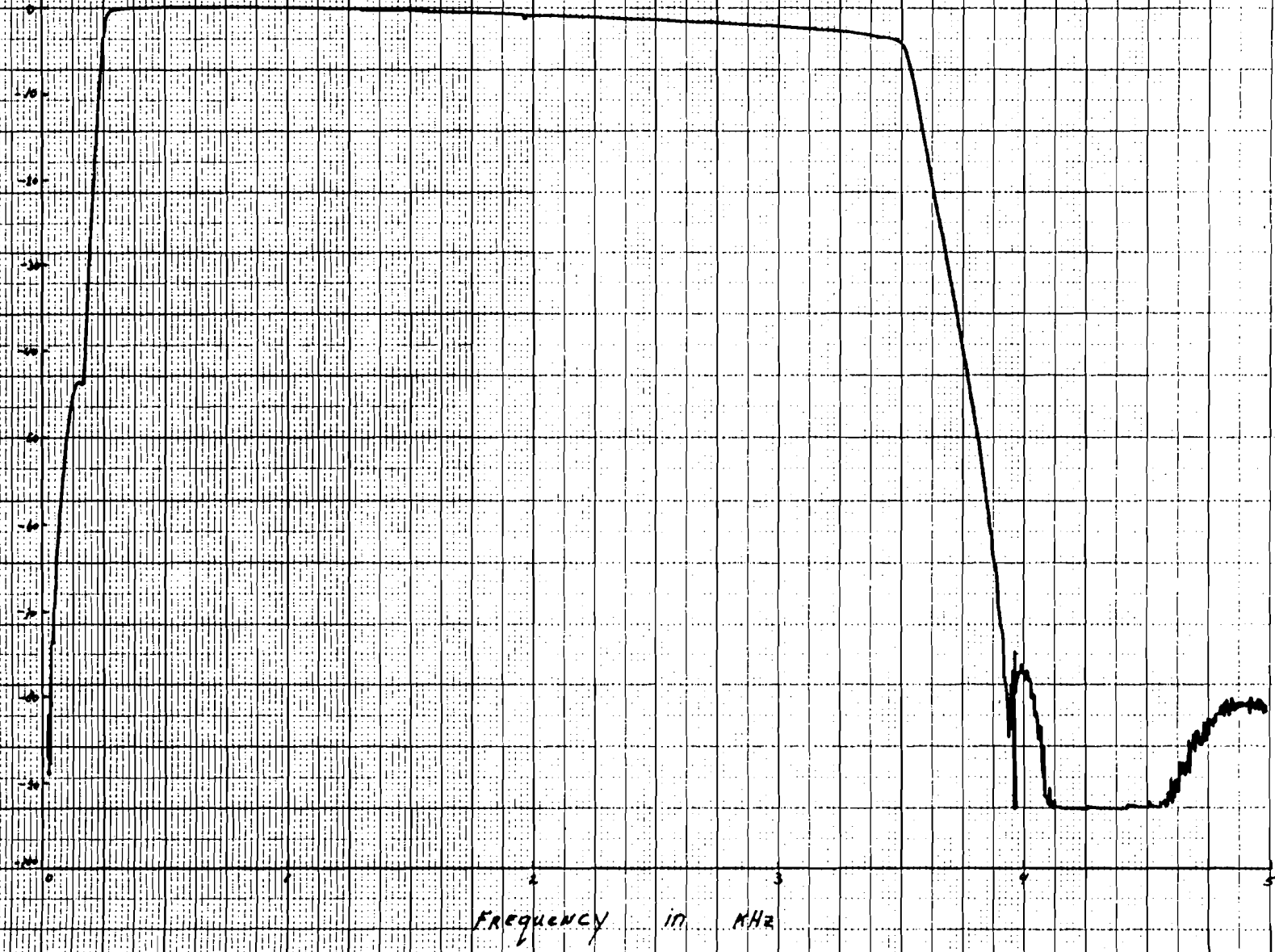
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APPENDIX A

Frequency response of bandlimiting filter

magnitude in decibels



APPENDIX B

Main program SBC transmitter part

```
*****
*                                     *
*                               CODER PROGRAM                               *
*                                     *
* THIS MAIN PROGRAM USES THE SUBROUTINES:                                *
*                                     *
*                               -QMF FILTERING                             *
*                               -PCMAQB CODING                            *
*                               -MULTIPLEXING                             *
*                                     *
* TO REALIZE A 16 KBIT PER SECOND SUBBAND CODER.                        *
* SYSTEM CLOCK RATE: 16 kHz                                             *
*-----*
* REPORTS FOR DETAILS:                                                 *
* PT REPORT BY P.CHITAMU                                               *
* FINAL REPORT BY J.BOLT                                               *
* FINAL REPORT BY R.KLEUTERS                                           *
*-----*
* AUTHOR: ROLAND KLEUTERS                                             *
* DATE: 29-5-1987                                                       *
*****
*
*-----*
* DATA MEMORY INITIALIZATION                                           *
*                                     default page=0                      *
*-----*
*
*.....*
*                               GENERAL VARIABLES DATA MEMORY PAGE 1    *
*.....*
*
CLOCK EQU 0          CLOCK RATE (16 kHz)
MODE EQU 1          MODE FOR ANALOG INTERFACE BOARD
AR00 EQU 2          STORAGE PLACE FOR AUXILIARY REGISTER 0
AR01 EQU 3          STORAGE PLACE FOR AUXILIARY REGISTER 1
ACH EQU 4           STORAGE PLACE FOR HIGH PART OF ACCUMULATOR
ACL EQU 5           STORAGE PLACE FOR LOW PART OF ACCUMULATOR
TREG EQU 6          STORAGE PLACE FOR T REGISTER
STATU EQU 7         STORAGE PLACE FOR STATUS WORD OF TMS32010
*                                     from now on default page
*.....*
*                               GENERAL VARIABLES                          *
*.....*
*
MSTAT EQU 0         MASK FOR STATE MOD 16
SWITCH EQU 1        COUNTS FIRST 4 INTERRUPTS; WHEN 0 MUX IS ON (00)
FRAME EQU 2         COUNTER FOR FRAME LENGTH (00)
ONE EQU 3           CHECK BIT (00)
STRAT EQU 4         CONTAINS 0, 1, or 2: VOICED, INTERM. or UNVO. (00)
SSTRAT EQU 5        BUFFERING (00)
STATE EQU 6         TREE POINTER (00)
RES0 EQU 7          INTERMEDIATE CALCULATION RESULTS FOR MAIN ROUTINE
RES1 EQU 8          INTERMEDIATE CALCULATION RESULTS FOR MAIN ROUTINE
PATH EQU 9          CONTAINS STATE/2=CHANNEL TO PROCESS (00)
*
```

```
*.....
*
*                FIRST QMF STAGE VARIABLES
*.....
*
*  BUFFER VARIABLES OF LOW1 (0-2 kHz) AND HIGH1 (2-4 kHz)          FORMAT=Q16
A1      EQU  10
B1      EQU  11
C1      EQU  12
*
*  INPUT TO (FILTER PART OF) FIRST QMF STAGE                      FORMAT=Q15
INPF    EQU  13
*
*  DELAY VARIABLES OF LOW1 (0-2 kHz)                               FORMAT=Q15
XF2     EQU  14          AFTER FILTERING INPUT STORED IN XF2, SO
XF3     EQU  15          XF1 NOT NEEDED
XF4     EQU  16
XF5     EQU  17
XF6     EQU  18
XF7     EQU  19
XF8     EQU  20
XF9     EQU  21
XF10    EQU  22
XF11    EQU  23
XF12    EQU  24
XF13    EQU  25
XF14    EQU  26
XF15    EQU  27
XF16    EQU  28
*
*  DELAY VARIABLES OF HIGH1 (2-4 kHz)                             FORMAT=Q15
XF18    EQU  29          AFTER FILTERING INPUT STORED IN XF18, SO
XF19    EQU  30          XF17 NOT NEEDED
XF20    EQU  31
XF21    EQU  32
XF22    EQU  33
XF23    EQU  34
XF24    EQU  35
XF25    EQU  36
XF26    EQU  37
XF27    EQU  38
XF28    EQU  39
XF29    EQU  40
XF30    EQU  41
XF31    EQU  42
XF32    EQU  43
*
*.....
*
*                SECOND QMF STAGE VARIABLES
*.....
*
*  BUFFER VARIABLES OF LOW2 (0-1 kHz) AND HIGH2 (1-2 kHz)        FORMAT=Q15
A2      EQU  44
B2      EQU  45
C2      EQU  46
*
```



```
* BUFFER VARIABLES OF LOW3 (3-4 kHz) AND HIGH3 (2-3 kHz)          FORMAT=Q17
A3   EQU  47
B3   EQU  48
C3   EQU  49
*
* INPUT TO (FILTER PART OF) SECOND QMF STAGE
INPS  EQU  50
*
* DELAY VARIABLES OF LOW2,HIGH2,LOW3 AND HIGH3 (TIME SHARING)
XS2  EQU  51
XS3  EQU  52
XS4  EQU  53
XS5  EQU  54
XS6  EQU  55
XS7  EQU  56
XS8  EQU  57
*
*.....
*                               THIRD QMF STAGE VARIABLES
*.....
* BUFFER VARIABLES OF LOW4 (0-500 Hz) AND HIGH4 (500-1000 Hz)    FORMAT=Q15
A4   EQU  58
B4   EQU  59
C4   EQU  60
*
* BUFFER VARIABLES OF LOW5 (1500-2000 Hz) AND HIGH5 (1000-1500 Hz) FOR.=Q15
A5   EQU  61
B5   EQU  62
C5   EQU  63
*
* BUFFER VARIABLES OF LOW7 (2000-2500 Hz) AND HIGH7 (2500-3000 Hz) FOR.=Q18
A7   EQU  64
B7   EQU  65
C7   EQU  66
*
* INPUT TO (FILTER PART OF) THIRD QMF STAGE
INPT  EQU  67
*
* DELAY VARIABLES OF LOW4,HIGH4,LOW5,HIGH5,LOW7 AND HIGH7 (TIME SHARING)
XT2  EQU  68
XT3  EQU  69
XT4  EQU  70
XT5  EQU  71
XT6  EQU  72
*
*.....
*                               VOICING DECISION VARIABLES
*.....
*
ENLOW EQU  73          ENERGY IN 0-2 kHz BAND (Q14)
ENHIGH EQU  74         ENERGY IN 2-4 kHz BAND (Q16)
*
```

*.....
* FIRST QMF STAGE COEFFICIENTS FORMAT=Q16
*.....

*
CF0 EQU 75
CF1 EQU 76
CF2 EQU 77
CF3 EQU 78
CF4 EQU 79
CF5 EQU 80
CF6 EQU 81
CF7 EQU 82
CF8 EQU 83
CF9 EQU 84
CF10 EQU 85
CF11 EQU 86
CF12 EQU 87
CF13 EQU 88
CF14 EQU 89
CF15 EQU 90

*.....
* SECOND QMF STAGE COEFFICIENTS FORMAT=Q16
*.....

*
CS0 EQU 91
CS1 EQU 92
CS2 EQU 93
CS3 EQU 94
CS4 EQU 95
CS5 EQU 96
CS6 EQU 97
CS7 EQU 98

*.....
* THIRD QMF STAGE COEFFICIENTS FORMAT=Q16
*.....

*
CT0 EQU 99
CT1 EQU 100
CT2 EQU 101
CT3 EQU 102
CT4 EQU 103
CT5 EQU 104

*.....
* PCMAQB VARIABLES
*.....

*
INPCM EQU 105 INPUT TO PCMAQB CODER
CALC1 EQU 106 GENERAL VARIABLE TO CALCULATE WITH
CALC2 EQU 107 GENERAL VARIABLE TO CALCULATE WITH
NBIT EQU 108 CONTAINS NR. OF BITS TO QUANTIZE TO (Q0)
X0 EQU 109 LOWEST POSSIBLE INPUT VALUE OF EXPONENT TABLE (Q14)
X127 EQU 110 HIGHEST POSSIBLE INPUT VALUE OF EXPONENT TABLE (Q14)
STEP EQU 111 1/(DISTANCE BETWEEN INPUTS OF EXPONENT TABLE) (Q7)
GAMMA EQU 112 GAMMA VALUE OF FORMULA (0.99) (Q15)
*

```
*.....
*                               MULTIPLEXER VARIABLES
*.....
```

```
*
COUNT EQU 113          COUNT FOR BITS TO SEND             (Q0)
BAND   EQU 114          FLAG, CURRENT SUBBAND PROCESSING    (Q0)
DATA   EQU 115          STORES OUTPUT BITS FOR SIGOUT
SIGOUT EQU 116          VARIABLE TO WRITE OUT SERIAL DATA FROM INSIDE
DEMO   EQU 117          NEW (FAW or SUBBAND) ACCESS MARK
*
-----
```

```
* PROGRAM MEMORY INITIALIZATION
*-----
```

```
*
ADR6 0
B START          BRANCH TO MAIN ROUTINE
B INTRO         BRANCH TO INTERRUPT SERVICE ROUTINE
ADR6 8          OTHERWISE ERRORS WITH TBLW'S
*
-----
```

```
*                               GENERAL CONSTANTS
*.....
```

```
*
CCLOCK DATA 312       CLOCK RATE (16 kHz)
MMODE DATA 24        MODE FOR ANALOG INTERFACE BOARD
MMSTAT DATA >000F    MASK FOR STATE MOD 16
*
-----
```

```
*                               SECOND QMF STAGE VARIABLES
*.....
```

```
* DELAY VARIABLES OF LOW2 (0-1 kHz)                               FORMAT=Q14
```

```
XXS2 DATA 0          AFTER FILTERING INPUT STORED IN XXS2, SO
XXS3 DATA 0          XXS1 NOT NEEDED
XXS4 DATA 0
XXS5 DATA 0
XXS6 DATA 0
XXS7 DATA 0
XXS8 DATA 0
*
-----
```

```
* DELAY VARIABLES OF HIGH2 (1-2 kHz)                               FORMAT=Q14
```

```
XXS10 DATA 0         AFTER FILTERING INPUT STORED IN XXS10, SO
XXS11 DATA 0         XXS9 NOT NEEDED
XXS12 DATA 0
XXS13 DATA 0
XXS14 DATA 0
XXS15 DATA 0
XXS16 DATA 0
*
-----
```

```
* DELAY VARIABLES OF LOW3 (3-4 kHz)                               FORMAT=Q16
```

```
XXS18 DATA 0         AFTER FILTERING INPUT STORED IN XXS18, SO
XXS19 DATA 0         XXS17 NOT NEEDED
XXS20 DATA 0
XXS21 DATA 0
XXS22 DATA 0
XXS23 DATA 0
XXS24 DATA 0
*
-----
```

```

* DELAY VARIABLES OF HIGH3 (2-3 kHz)                          FORMAT=Q16
XXS26 DATA 0 AFTER FILTERING INPUT STORED IN XXS26, SO
XXS27 DATA 0 XXS25 NOT NEEDED
XXS28 DATA 0
XXS29 DATA 0
XXS30 DATA 0
XXS31 DATA 0
XXS32 DATA 0

```

```

*
*.....
* THIRD QMF STAGE VARIABLES
*.....

```

```

* DELAY VARIABLES OF LOW4 (0-500 Hz)                          FORMAT=Q14
XXT2 DATA 0 AFTER FILTERING INPUT STORED IN XXT2, SO
XXT3 DATA 0 XXT1 NOT NEEDED
XXT4 DATA 0
XXT5 DATA 0
XXT6 DATA 0

```

```

* DELAY VARIABLES OF HIGH4 (500-1000 Hz)                      FORMAT=Q14
XXT8 DATA 0 AFTER FILTERING INPUT STORED IN XXT8, SO
XXT9 DATA 0 XXT7 NOT NEEDED
XXT10 DATA 0
XXT11 DATA 0
XXT12 DATA 0

```

```

* DELAY VARIABLES OF LOW5 (1500-2000 Hz)                      FORMAT=Q14
XXT14 DATA 0 AFTER FILTERING INPUT STORED IN XXT14, SO
XXT15 DATA 0 XXT13 NOT NEEDED
XXT16 DATA 0
XXT17 DATA 0
XXT18 DATA 0

```

```

* DELAY VARIABLES OF HIGH5 (1000-1500 Hz)                     FORMAT=Q14
XXT20 DATA 0 AFTER FILTERING INPUT STORED IN XXT20, SO
XXT21 DATA 0 XXT19 NOT NEEDED
XXT22 DATA 0
XXT23 DATA 0
XXT24 DATA 0

```

```

* DELAY VARIABLES OF LOW7 (2000-2500 Hz)                       FORMAT=Q17
XXT26 DATA 0 AFTER FILTERING INPUT STORED IN XXT26, SO
XXT27 DATA 0 XXT25 NOT NEEDED
XXT28 DATA 0
XXT29 DATA 0
XXT30 DATA 0

```

```

* DELAY VARIABLES OF HIGH7 (2500-3000 Hz)                     FORMAT=Q17
XXT32 DATA 0 AFTER FILTERING INPUT STORED IN XXT32, SO
XXT33 DATA 0 XXT31 NOT NEEDED
XXT34 DATA 0
XXT35 DATA 0
XXT36 DATA 0

```

*

*.....
* FIRST QMF STAGE COEFFICIENTS FORMAT=Q16
*.....

*
CCF0 DATA 45
CCF1 DATA -92
CCF2 DATA -83
CCF3 DATA 277
CCF4 DATA 93
CCF5 DATA -620
CCF6 DATA -9
CCF7 DATA 1178
CCF8 DATA -274
CCF9 DATA -2047
CCF10 DATA 955
CCF11 DATA 3470
CCF12 DATA -2579
CCF13 DATA -6541
CCF14 DATA 8425
CCF15 DATA 30566

*.....
* SECOND QMF STAGE COEFFICIENTS FORMAT=Q16
*.....

*
CCS0 DATA 69
CCS1 DATA -331
CCS2 DATA -170
CCS3 DATA 1812
CCS4 DATA -633
CCS5 DATA -5924
CCS6 DATA 6409
CCS7 DATA 31525

*.....
* THIRD QMF STAGE COEFFICIENTS FORMAT=Q16
*.....

*
CCT0 DATA -250
CCT1 DATA 1236
CCT2 DATA -178
CCT3 DATA -5551
CCT4 DATA 5798
CCT5 DATA 31745

*.....
* PCMAQB'S TABLES
*.....

* (1-GAMMA)*DC FOR 6 SUBBANDS FORMAT=Q22
0GD1 DATA -10486 SUBBAND 1
DUM1 DATA 0
0GD4 DATA -13642 SUBBAND 4
0GD5 DATA -23112 SUBBAND 5
0GD2 DATA -10486 SUBBAND 2
DUM2 DATA 0
0GD3 DATA -13642 SUBBAND 3
0GD6 DATA -23112 SUBBAND 6

```
*
*      D(N) FOR 6 SUBBANDS                                FORMAT=Q14
DN1   DATA -13563      SUBBAND 1
DUM3  DATA  0
DN4   DATA -14796      SUBBAND 4
DN5   DATA -18495      SUBBAND 5
DN2   DATA -13563      SUBBAND 2
DUM4  DATA  0
DN3   DATA -14796      SUBBAND 3
DN6   DATA -18495      SUBBAND 6
*
*      BIT ALLOCATION FOR VOICED STRATEGY (0)              FORMAT=Q0
BAVOI1 DATA 4          SUBBAND 1
DUM5   DATA 0
BAVOI4 DATA 2          SUBBAND 4
BAVOI5 DATA 2          SUBBAND 5
BAVOI2 DATA 4          SUBBAND 2
DUM6   DATA 0
BAVOI3 DATA 3          SUBBAND 3
BAVOI6 DATA 0          SUBBAND 6
*
*      BIT ALLOCATION FOR INTERMEDIATE STRATEGY (1)        FORMAT=Q0
BAINT1 DATA 4          SUBBAND 1
DUM7   DATA 0
BAINT4 DATA 2          SUBBAND 4
BAINT5 DATA 2          SUBBAND 5
BAINT2 DATA 3          SUBBAND 2
DUM8   DATA 0
BAINT3 DATA 2          SUBBAND 3
BAINT6 DATA 2          SUBBAND 6
*
*      BIT ALLOCATION FOR UNVOICED STRATEGY (2)            FORMAT=Q0
BAUNV1 DATA 2          SUBBAND 1
DUM9   DATA 0
BAUNV4 DATA 3          SUBBAND 4
BAUNV5 DATA 2          SUBBAND 5
BAUNV2 DATA 3          SUBBAND 2
DUM10  DATA 0
BAUNV3 DATA 3          SUBBAND 3
BAUNV6 DATA 2          SUBBAND 6
*
*      STORAGE OF CODED SUBBAND SAMPLES                   FORMAT=Q0
COD1   DATA 0          SUBBAND 1
DUM11  DATA 0
COD4   DATA 0          SUBBAND 4
COD5   DATA 0          SUBBAND 5
COD2   DATA 0          SUBBAND 2
DUM12  DATA 0
COD3   DATA 0          SUBBAND 3
COD6   DATA 0          SUBBAND 6
*
*      SATURATION LEVELS FOR QUANTIZER                    FORMAT=Q0
SAT1   DATA 1          2 BITS
SAT2   DATA 3          3 BITS
SAT3   DATA 7          4 BITS
*
```

* QUANTIZER CONSTANTS

SSTEP DATA 19637 1/(DISTANCE BETWEEN INPUTS OF EXPONENT TABLE) (Q7)
GAMMA DATA 32440 GAMMA=0.99 (Q15)

*

* LOGTABLE 2 BIT CASE

FORMAT=Q18

LOG20 DATA 18268
LOG21 DATA -4626
LOG22 DATA -4626
LOG23 DATA 18268
DUM13 DATA 0
DUM14 DATA 0
DUM15 DATA 0
DUM16 DATA 0
DUM17 DATA 0
DUM18 DATA 0
DUM19 DATA 0
DUM20 DATA 0
DUM21 DATA 0
DUM22 DATA 0

*

* LOGTABLE 3 BIT CASE

FORMAT=Q18

LOG30 DATA 11540
LOG31 DATA 0
LOG32 DATA 0
LOG33 DATA -4626
LOG34 DATA -4626
LOG35 DATA 0
LOG36 DATA 0
LOG37 DATA 11540
DUM23 DATA 0
DUM24 DATA 0
DUM25 DATA 0
DUM26 DATA 0

*

* LOGTABLE 4 BIT CASE

FORMAT=Q18

LOG40 DATA 24918
LOG41 DATA 19728
LOG42 DATA 13377
LOG43 DATA 5189
LOG44 DATA -2999
LOG45 DATA -2999
LOG46 DATA -2999
LOG47 DATA -2999
LOG48 DATA -2999
LOG49 DATA -2999
LOG410 DATA -2999
LOG411 DATA -2999
LOG412 DATA 5189
LOG413 DATA 13377
LOG414 DATA 19728
LOG415 DATA 24918

*

* LOWEST AND HIGHEST POSSIBLE INPUTS OF EXPONENT TABLE FORMAT=Q14
XX10 DATA -13563 SUBBAND 1
XX1127 DATA 0 SUBBAND 1
DUM27 DATA 0
DUM28 DATA 0
XX40 DATA -14796 SUBBAND 4
XX4127 DATA -1233 SUBBAND 4
XX50 DATA -18495 SUBBAND 5
XX5127 DATA -4932 SUBBAND 5
XX20 DATA -13563 SUBBAND 2
XX2127 DATA 0 SUBBAND 2
DUM29 DATA 0
DUM30 DATA 0
XX30 DATA -14796 SUBBAND 3
XX3127 DATA -1233 SUBBAND 3
XX60 DATA -18495 SUBBAND 6
XX6127 DATA -4932 SUBBAND 6

*
* EXPONENT TABLE OF CODER
* FORMATS ARE: SUBBAND 1 Q4
* SUBBAND 2 Q4
* SUBBAND 3 Q3
* SUBBAND 4 Q3
* SUBBAND 5 Q0
* SUBBAND 6 Q0
*

CEXPO DATA 32767
 DATA 30859
 DATA 29060
 DATA 27367
 DATA 25772
 DATA 24271
 DATA 22856
 DATA 21525
 DATA 20270
 DATA 19089
 DATA 17977
 DATA 16929
 DATA 15943
 DATA 15014
 DATA 14139
 DATA 13315
 DATA 12539
 DATA 11809
 DATA 11121
 DATA 10473
 DATA 9862
 DATA 9288
 DATA 8746
 DATA 8237
 DATA 7757
 DATA 7305
 DATA 6879
 DATA 6478
 DATA 6101
 DATA 5745
 DATA 5411
 DATA 5095

DATA 4798
DATA 4519
DATA 4256
DATA 4008
DATA 3774
DATA 3554
DATA 3347
DATA 3152
DATA 2968
DATA 2795
DATA 2632
DATA 2479
DATA 2335
DATA 2199
DATA 2070
DATA 1950
DATA 1836
DATA 1729
DATA 1628
DATA 1534
DATA 1444
DATA 1360
DATA 1281
DATA 1206
DATA 1136
DATA 1070
DATA 1007
DATA 949
DATA 893
DATA 841
DATA 792
DATA 746
DATA 703
DATA 662
DATA 623
DATA 587
DATA 553
DATA 520
DATA 490
DATA 462
DATA 435
DATA 409
DATA 385
DATA 363
DATA 342
DATA 322
DATA 303
DATA 286
DATA 269
DATA 253
DATA 238
DATA 225
DATA 211
DATA 199
DATA 188
DATA 177
DATA 166
DATA 157

DATA 148
DATA 139
DATA 131
DATA 123
DATA 116
DATA 109
DATA 103
DATA 97
DATA 91
DATA 86
DATA 81
DATA 76
DATA 72
DATA 68
DATA 64
DATA 60
DATA 56
DATA 53
DATA 50
DATA 47
DATA 44
DATA 42
DATA 39
DATA 37
DATA 35
DATA 33
DATA 31
DATA 29
DATA 27
DATA 26
DATA 24
DATA 23
DATA 22
DATA 20
DATA 19
DATA 18
DATA 17
DATA 16

*

* MAIN ROUTINE

*

*

.....
* INITIALIZATIONS
.....

*

START DINT

*

* INITIALIZE STATUS BITS

LARP 0

LDPK 0

SOVM

*

```
* INITIALIZE GENERAL VARIABLES OF DATA MEMORY PAGE 1
  LDPK 1
  LACK CCLOCK
  TBLR CLOCK          CLOCK RATE (16 kHz)
  LACK MMODE
  TBLR MODE           MODE FOR ANALOG INTERFACE BOARD
  LDPK 0
*
* INITIALIZE GENERAL VARIABLES OF DATA MEMORY PAGE 0
  LACK MMSTAT
  TBLR MSTAT          MASK FOR STATE MOD 16
  LACK 5              NO. OF WAIT CYCLI BEFOR MULTIPLEXER START
  SACL SWITCH         COUNTER FOR MULTIPLEXER ON or OFF, NOW OFF
  ZAC
  SACL FRAME          STATUS FOR CODER
  LACK 1
  SACL ONE            '1' FOR LOGICAL MANIPULATIONS
  SACL STRAT          BY DEFAULT LOAD INTERMEDIATE STRATEGY
  SACL SSTRAT         ALSO INTERMEDIATE STRATEGY FOR BUFFER
  ZAC
  SACL STATE          INIT TREE POINTER
*
* INITIALIZE PCMAQR VARIABLES
  LACK SSTEP
  TBLR STEP           1/(DISTANCE BETWEEN INPUTS OF EXPONENT TABLE)
  LACK 6GAMMA
  TBLR GAMMA          GAMMA VALUE OF FORMULA (0.99)
*
* INITIALIZE MULTIPLEXER VARIABLES
  ZAC
  SACL BAND           STATUS FOR MULTIPLEXER
  LACK >FF
  SACL DEMO           NEW ACCESS MARK
*
* LOAD QMF COEFFICIENTS INTO DATA RAM
  LARK 1,CT5
  LARK 0,29
  LACK CCT5
LOAD  LARP 1
      TBLR *-,0
      SUB ONE
      BANZ LOAD
*
* INITIALIZE ANALOG INTERFACE BOARD (AIB)
  LDPK 1
  OUT MODE,0          PROGRAM AIB
  OUT CLOCK,1         SET CLOCK RATE OF AIB
  LDPK 0
  OUT RES0,2          CLEAR AD1S (PROTECTION AGAINST RESET INTERRUPTS)
  OUT RES0,3          CLEAR AD2S (PROTECTION AGAINST RESET INTERRUPTS)
*
* INITIALIZATION AND VARIABLE LOADING DONE, NOW BEGIN
  EINT
*
```

```
* DISCARD FIRST TWO INTERRUPTS TO GET A DEFINED STARTING POINT
ACKNO ZALS STATE
      BZ ACKNO
      LACK 1
      XOR STATE
      BZ ACKNO
      ZAC
      SACL STATE          INIT TREE PGINTER
*
* .....
*                               AORTA
* .....
*
* ODD CYCLE OF 16 kHz CLOCK MUST JUST HAVE PASSED
WAITEV LAC STATE
      AND ONE
      BNZ WAITEV
* EVEN CYCLE OF 16 kHz CLOCK HAS PASSED (BIT 0 OF STATE=0)
WAIT LAC STATE
      AND ONE
      BZ WAIT
* ODD CYCLE OF 16 kHz CLOCK HAS JUST PASSED (BIT 0 OF STATE=1)
*
* INPUT SAMPLE RATE REDUCED FROM 16 kHz (2*8 kHz) TO 8 kHz BY OMITTING
* EVERY EVEN INPUT SAMPLE
*
* DURING THE PROCESSING OF A SUBBAND SAMPLE STATE CAN CHANGE AS A
* RESULT OF AN INTERRUPT, SO DEFINE PATH=STATE/2
      LAC STATE,15
      SACH PATH
*
* ACTUAL SUBBAND CODING
      CALL QMF
      CALL PCMAQB
      B WAITEV
*
*-----
* INTERRUPT SERVICE ROUTINE
*-----
*
INTRO DINT
*
* SAVE STATUS OF TMS32010
      SST STATU
      MPY ONE          T REGISTER TO P REGISTER
      LDPK 1
      SAR 0,AR00
      SAR 1,AR01
      SACH ACH
      SACL ACL
      PAC
      SACL TREG
      LDPK 0
*
```

```
* STATE UPDATE
  LAC STATE
  ADD ONE
  AND MSTAT          STATE MOD 16
  SACL STATE
*
* SAMPLE INPUT (16 kHz = 2*8 kHz)
  IN INPF,2          ONE OF TWO SAMPLES IS USED LATER
*
* CHECK IF MULTIPLEXER IS ON (SWITCH=0); IF NOT DECREMENT SWITCH
* AND RETURN TO PLACE OF INTERRUPT
  ZALS SWITCH
  BZ MUXER
  SUB ONE
  SACL SWITCH
  B RSSTAT
*
* RUN THROUGH MULTIPLEXER
MUXER CALL MUX
*
* RESTORE STATUS OF TMS32010
RSSTAT LDPK 1
      LAR 0,AR00
      LAR 1,AR01
      ZALH ACH
      ADDS ACL
      LT TREG
      LST STATU
*
* RETURN TO PLACE OF INTERRUPT AND CONTINUE
  EINT
  RET
```

CODER PROGRAM DONE

=====

APPENDIX C
Subroutine QMF

```
*****
#                                     QMF FILTERING                                     #
#                                                                                     #
# THIS SUBROUTINE INCORPORATES:                                                       #
#     -QMF FILTER TREE                                                                #
#     -ENERGY MEASUREMENT                                                             #
#     -VOICING STRATEGY DECISION                                                      #
#     -TAKE OVER OF NEW VOICING STRATEGY                                             #
#     FOR THE PCM CODING                                                              #
#                                                                                     #
# INPUT: 8 kHz SAMPLES (1 STREAM)                                                    #
# OUTPUT: FAW + 1 kHz SAMPLES (8 STREAMS)                                           #
#-----#
# SUBBAND PROCESSING : 1,8,4,5,2,7,3,6                                             #
#                                                                                     #
# BIT ALLOCATION :                                                                    #
#           ----- 1  2  3  4  5  6  7  8 ----- #
#           VOICED   4  4  3  2  2  0  0  0 #
#           INTERMEDIATE 4  3  2  2  2  2  0  0 #
#           UNVOICED  2  3  3  3  2  2  0  0 #
# FAW's :                                                                              #
#           VOICED >76      INTERMEDIATE >96      UNVOICED >CC #
#-----#
# AUTHOR: ROLAND KLEUTERS                                                            #
# DATE: 29-5-1987                                                                    #
#-----#
#                                                                                     #
#-----#
# MACRO'S                                                                            #
#-----#
#                                                                                     #
#     MACRO TO REALIZE FILTERING IN LOW BRANCH OF FIRST QMF STAGE                   #
FLTFL0 $MACRO X
      ZAC
      LT  XF16
      MPY CF0
      LTD XF15
      MPY CF2
      LTD XF14
      MPY CF4
      LTD XF13
      MPY CF6
      LTD XF12
      MPY CF8
      LTD XF11
      MPY CF10
      LTD XF10
      MPY CF12
      LTD XF9
      MPY CF14
      LTD XF8
      MPY CF15
      LTD XF7
      MPY CF13
      LTD XF6
      MPY CF11
      LTD XF5
      MPY CF9
```

```
LTD XF4
MPY CF7
LTD XF3
MPY CF5
LTD XF2
MPY CF3
LTA INPF      8 kHz INPUT TO FIRST QMF STAGE
MPY CF1
APAC
SACH :X.S:,1
LAC INPF
SACL XF2
$END
```

*

* MACRO TO REALIZE FILTERING IN HIGH BRANCH OF FIRST QMF STAGE

```
FLTFHI $MACRO X
ZAC
LT XF32
MPY CF1
LTD XF31
MPY CF3
LTD XF30
MPY CF5
LTD XF29
MPY CF7
LTD XF28
MPY CF9
LTD XF27
MPY CF11
LTD XF26
MPY CF13
LTD XF25
MPY CF15
LTD XF24
MPY CF14
LTD XF23
MPY CF12
LTD XF22
MPY CF10
LTD XF21
MPY CF8
LTD XF20
MPY CF6
LTD XF19
MPY CF4
LTD XF18
MPY CF2
LTA INPF      8 kHz INPUT TO FIRST QMF STAGE
MPY CF0
APAC
SACH :X.S:,1
LAC INPF
SACL XF18
$END
```

*

* MACRO TO REALIZE FILTERING IN LOW BRANCH OF SECOND QMF STAGE

```
FLTSLO $MACRO X
      ZAC
      LT  XS8
      MPY CS0
      LTD XS7
      MPY CS2
      LTD XS6
      MPY CS4
      LTD XS5
      MPY CS6
      LTD XS4
      MPY CS7
      LTD XS3
      MPY CS5
      LTD XS2
      MPY CS3
      LTD IMPS          4 kHz INPUT TO SECOND QMF STAGE
      MPY CS1
      APAC
      SACH :X.S:,1
      $END
```

* MACRO TO REALIZE FILTERING IN HIGH BRANCH OF SECOND QMF STAGE

```
FLTSHI $MACRO X
      ZAC
      LT  XS8
      MPY CS1
      LTD XS7
      MPY CS3
      LTD XS6
      MPY CS5
      LTD XS5
      MPY CS7
      LTD XS4
      MPY CS6
      LTD XS3
      MPY CS4
      LTD XS2
      MPY CS2
      LTD IMPS          4 kHz INPUT TO SECOND QMF STAGE
      MPY CS0
      APAC
      SACH :X.S:,1
      $END
```

* MACRO TO REALIZE FILTERING IN LOW BRANCH OF THIRD QMF STAGE

```
FLTTLO $MACRO X
      ZAC
      LT  XT6
      MPY CT0
      LTD XT5
      MPY CT2
      LTD XT4
      MPY CT4
      LTD XT3
      MPY CT5
      LTD XT2
```

```
MPY CT3
LTD INPT          2 kHz INPUT TO THIRD QMF STAGE
MPY CT1
APAC
SACH :X.S:,1
$END
```

```
*
* MACRO TO REALIZE FILTERING IN HIGH BRANCH OF THIRD QMF STAGE
FLTTHI $MACRO X
ZAC
LT XT6
MPY CT1
LTD XT5
MPY CT3
LTD XT4
MPY CT5
LTD XT3
MPY CT4
LTD XT2
MPY CT2
LTD INPT          2 kHz INPUT TO THIRD QMF STAGE
MPY CT0
APAC
SACH :X.S:,1
$END
```

```
*
* MACRO TO LOAD DELAY VARIABLES OF SECOND QMF STAGE
LOADDS $MACRO X
LACK :X.S:
LARK 0,XS2
TBLR ++
ADD ONE
TBLR ++
ADD ONE
TBLR ++
ADD ONE
TBLR ++
ADD ONE
TBLR ++
ADD ONE
TBLR ++
ADD ONE
TBLR +
$END
```

```
*
* MACRO TO SAVE DELAY VARIABLES OF SECOND QMF STAGE
SAVEDS $MACRO X
LACK :X.S:
LARK 0,XS2
TBLW ++
ADD ONE
TBLW ++
ADD ONE
TBLW ++
ADD ONE
TBLW ++
ADD ONE
TBLW ++
$END
```

```
ADD ONE
TBLW **
ADD ONE
TBLW *
$END

*
* MACRO TO LOAD DELAY VARIABLES OF THIRD QMF STAGE
LOADDT $MACRO X
LACK :X,S:
LARK 0,XT2
TBLR **
ADD ONE
TBLR **
ADD ONE
TBLR **
ADD ONE
TBLR **
ADD ONE
TBLR **
ADD ONE
TBLR *
$END

*
* MACRO TO SAVE DELAY VARIABLES OF THIRD QMF STAGE
SAVEDT $MACRO X
LACK :X,S:
LARK 0,XT2
TBLW **
ADD ONE
TBLW **
ADD ONE
TBLW **
ADD ONE
TBLW **
ADD ONE
TBLW **
ADD ONE
TBLW *
$END

*
*-----*
* MACRO DEFINITION DONE; MAIN PART OF SUBROUTINE BEGINS
*-----*
*
QMF      NOP                ENTRY POINT
*
* SELECT LOW1 OR HIGH1 FOR FIRST QMF STAGE
LAC ONE          BIT 0=0 or 1
AND PATH
BZ LOW1
B HIGH1
*
*-----*
* FIRST QMF STAGE:
*          LOW1 0kHz-2kHz
*          HIGH1 2kHz-4kHz
*-----*
*
LOW1      DMOV A1          DOUBLE BUFFER
*
```

```
* OUTPUT CALCULATION
  LAC C1,14
  ADD B1,14
  SACH INPS          4 kHz INPUT TO SECOND QMF STAGE
*
* FILTER 0-2 kHz
  FLTFO A1
*
* ENERGY SUMMATION, SUM X(n) SQUARED
  LT INPS
  MPY INPS
  PAC              ACCU CONTAINS INPS*INPS IN Q28 FORMAT
  ADD ENLOW,14    ACCU CONTAINS ENLOW IN Q28 FORMAT
  SACL RES0      MULTIPLY ACCU WITH 4
  SACH RES1
  ADD RES0
  ADDH RES1
  ADD RES0
  ADDH RES1
  ADD RES0
  ADDH RES1      ACCU CONTAINS ENLOW IN Q30 FORMAT
  SACH ENLOW     ENLOW IN Q14 FORMAT
*
* SELECT LOW2 OR HIGH2 FOR SECOND QMF STAGE
  LAC ONE,1
  AND PATH
  BZ LOW2
  B HIGH2
*
HIGH1  NOP          NO DOUBLE BUFFER
*
* OUTPUT CALCULATION
  ZALH C1
  SUBH B1
  SACH INPS          4 kHz INPUT TO SECOND QMF STAGE
*
* FILTER 2-4 kHz
  FLTFO C1
*
* ENERGY SUMMATION, SUM X(n) SQUARED
  LT INPS
  MPY INPS
  PAC              ACCU CONTAINS INPS*INPS IN Q32 FORMAT
  ADDH ENHIGH     ACCU CONTAINS ENHIGH IN Q32 FORMAT
  SACH ENHIGH     ENHIGH IN Q16 FORMAT
*
* SELECT LOW3 OR HIGH3 FOR SECOND QMF STAGE
  LAC ONE,1
  AND PATH
  BZ LOW3
  B HIGH3
*
```

* SECOND QMF STAGE:

* LOW2 0kHz-1kHz
* HIGH2 1kHz-2kHz
* LOW3 3kHz-4kHz
* HIGH3 2kHz-3kHz

*
* LOW2 DMOV A2 DOUBLE BUFFER
*
* OUTPUT CALCULATION
* LAC C2,15
* ADD B2,15
* SACH INPT 2 kHz INPUT TO THIRD QMF STAGE
*
* LOAD DELAY VARIABLES OF LOW2
* LOADDS XIXS2
*
* FILTER 0-1 kHz
* FLTSLO A2
*
* SAVE DELAY VARIABLES OF LOW2
* SAVEDS XIXS2
*
* SELECT LOW4 OR HIGH4 FOR THIRD QMF STAGE
* LAC ONE,2
* AND PATH
* BZ LOW4
* B HIGH4
*
* HIGH2 NOP NO DOUBLE BUFFER
*
* OUTPUT CALCULATION
* LAC C2,15
* SUB B2,15
* SACH INPT 2 kHz INPUT TO THIRD QMF STAGE
*
* LOAD DELAY VARIABLES OF HIGH2
* LOADDS XIXS10
*
* FILTER 1-2 kHz
* FLTSHI C2
*
* SAVE DELAY VARIABLES OF HIGH2
* SAVEDS XIXS10
*
* SELECT LOW5 OR HIGH5 FOR THIRD QMF STAGE
* LAC ONE,2
* AND PATH
* BZ LOW5
* B HIGH5
*
* LOW3 DMOV A3 DOUBLE BUFFER
*
* OUTPUT CALCULATION
* ZALM C3
* ADDDH B3
* SACH INPT 2 kHz INPUT TO THIRD QMF STAGE

```
*
* LOAD DELAY VARIABLES OF LOW3
*   LOADDS XYS18
*
* FILTER 3-4 kHz
*   FLTSLO A3
*
* SAVE DELAY VARIABLES OF LOW3
*   SAVEDS XYS18
*
* THE 3-4 kHz SUBBAND IS NOT CODED, SO FURTHER SPLITTING IS NOT NECESSARY
*
* VOICING DECISION?
* FRAME HAS TO BE 3 AND PATH HAS TO BE 5; INTERRUPT 75
*   LACK 3
*   XOR FRAME
*   BNZ NOVODE      FRAME=3?
*   LACK 5
*   XOR PATH       FRAME=5?
*   BZ VOIDEC
NOVODE RET          NOT INTERRUPT 75; NO VOICING DECISION; QMF DONE
*
HIGH3 NOP          NO DOUBLE BUFFER
*
* OUTPUT CALCULATION
*   ZALH C3
*   SUBH B3
*   SACH INPT      2 kHz INPUT TO THIRD QMF STAGE
*
* LOAD DELAY VARIABLES OF HIGH3
*   LOADDS XYS26
*
* FILTER 2-3 kHz
*   FLTSHI C3
*
* SAVE DELAY VARIABLES OF HIGH3
*   SAVEDS XYS26
*
* SELECT LOW7 OR HIGH7 FOR THIRD QMF STAGE
*   LAC ONE,2
*   AND PATH
*   BZ LOW7
*   B   HIGH7
*
*-----*
* THIRD QMF STAGE:
*   LOW4 0-500 Hz
*   HIGH4 500-1000 Hz
*   LOW5 1500-2000 Hz
*   HIGH5 1000-1500 Hz
*   LOW7 2000-2500 Hz
*   HIGH7 2500-3000 Hz
*   MAKE A VOICING DECISION
*   TAKE OVER NEW VOICING STRATEGY
*-----*
*
LOW4 DMOV A4      DOUBLE BUFFER
*
```

```
* OUTPUT CALCULATION
  LAC C4,15
  ADD B4,15
  SACH INPCM          1 kHz INPUT TO PCM CODER
*
* LOAD DELAY VARIABLES OF LOW4
  LOADDT XXT2
*
* FILTER 0-500 Hz
  FLTTLO A4
*
* SAVE DELAY VARIABLES OF LOW4
  SAVEDT XXT2
*
* QMF DONE;
* BEFORE EXIT, UPDATE FRAME AND TAKE OVER NEW STRATEGY IF END OF FRAME
  LAR 0,FRAME
  BANZ SFRAME
  LACK 7              END OF FRAME; PREPARE NEXT FRAME
  SACL FRAME
  DMOV STRAT         TAKE OVER NEW VOICING STRATEGY
  RET
SFRAME SAR 0,FRAME  NOT END OF FRAME; PREPARE NEXT DATABLOCK
  RET
*
HIGH4  NOP          NO DOUBLE BUFFER
*
* OUTPUT CALCULATION
  LAC C4,15
  SUB B4,15
  SACH INPCM          1 kHz INPUT TO PCM CODER
*
* LOAD DELAY VARIABLES OF HIGH4
  LOADDT XXT8
*
* FILTER 500-1000 Hz
  FLTTHI C4
*
* SAVE DELAY VARIABLES OF HIGH4
  SAVEDT XXT8
*
* QMF DONE
  RET
*
LOW5   DMOV A5          DOUBLE BUFFER
*
* OUTPUT CALCULATION
  ZALH C5
  ADDH B5
  SACH INPCM          1 kHz INPUT TO PCM CODER
*
* LOAD DELAY VARIABLES OF LOW5
  LOADDT XXT14
*
* FILTER 1500-2000 Hz
  FLTTLO A5
*
```

```
*   SAVE DELAY VARIABLES OF LOW5
      SAVEDT XXT14
*
*   QMF DONE
      RET
*
HIGH5  NOP                NO DOUBLE BUFFER
*
*   OUTPUT CALCULATION
      ZALH C5
      SUBH B5
      SACH INPCM          1 kHz INPUT TO PCM CODER
*
*   LOAD DELAY VARIABLES OF HIGH5
      LOADDT XXT20
*
*   FILTER 1000-1500 Hz
      FLTHI C5
*
*   SAVE DELAY VARIABLES OF HIGH5
      SAVEDT XXT20
*
*   QMF DONE
      RET
*
LOW7   DMOV A7            DOUBLE BUFFER
*
*   OUTPUT CALCULATION
      ZALH C7
      ADDH B7
      SACH INPCM          1 kHz INPUT TO PCM CODER
*
*   LOAD DELAY VARIABLES OF LOW7
      LOADDT XXT26
*
*   FILTER 2000-2500 Hz
      FLTLD A7
*
*   SAVE DELAY VARIABLES OF LOW7
      SAVEDT XXT26
*
*   QMF DONE
      RET
*
HIGH7  NOP                NO DOUBLE BUFFER
*
*   OUTPUT CALCULATION
      ZALH C7
      SUBH B7
      SACH INPCM          1 kHz INPUT TO PCM CODER
*
*   LOAD DELAY VARIABLES OF HIGH7
      LOADDT XXT32
*
*   FILTER 2500-3000 Hz
      FLTHI C7
*
```



```
*   SAVE DELAY VARIABLES OF HIGH7
      SAVEDT XXT32
*
*   QMF DONE
      RET
*
*   MAKE A VOICING DECISION
VOIDEC  NOP
*
*   IF 20ENHIGH-ENLOW LESS THAN OR EQUAL TO ZERO THEN VOICED
      LT  ENHIGH
      MPYK +20
      PAC
      SUB  ENLOW,2
      BLEZ SETV01
*
*   IF 3ENHIGH-2ENLOW GREATER THAN OR EQUAL TO ZERO THEN UNVOICED
      LT  ENHIGH
      MPYK +3
      PAC
      SUB  ENLOW,2
      SUB  ENLOW,2
      BGEZ SETUN
*
*   IF 10ENHIGH-ENLOW LESS THAN OR EQUAL TO ZERO AND PREVIOUS VOICING
*   WAS VOICED THEN VOICED
      LAC  STRAT
      BNZ  PREUNV
      LT  ENHIGH
      MPYK +10
      PAC
      SUB  ENLOW,2
      BLEZ SETV01
*
*   IF 3ENHIGH-ENLOW GREATER THAN OR EQUAL TO ZERO AND PREVIOUS VOICING
*   WAS UNVOICED THEN UNVOICED
PREUNV  LACK 2
      XOR  STRAT
      BNZ  SETINT
      LT  ENHIGH
      MPYK 3
      PAC
      SUB  ENLOW,2
      BGEZ SETUN
*
*   ELSE NEW STRATEGY IS INTERMEDIATE STRATEGY
*
*   SET INTERMEDIATE STRATEGY AND CLEAR VOICING DECISION BUFFERS
SETINT  LACK 1
      SACL STRAT
      ZAC
      SACL ENLOW
      SACL ENHIGH
      RET
      QMF DONE
*
```

* SET VOICED STRATEGY AND CLEAR VOICING DECISION BUFFERS

SETVOI ZAC
SACL STRAT
SACL ENLOW
SACL ENHIGH
RET QMF DONE

* SET UNVOICED STRATEGY AND CLEAR VOICING DECISION BUFFERS

SETUN LACK 2
SACL STRAT
ZAC
SACL ENLOW
SACL ENHIGH
RET QMF DONE

*
* QMF FILTERING DDNE

=====

APPENDIX D

Subroutine PCMAQB (coder)

```

*****
*                                     PCMA8B CODING                                     *
*                                     *
* THIS SUBROUTINE INCORPORATES :
*                                     -BACKWARD STEPSIZE ADAPTION OF THE QUANTIZER
*                                     -QUANTIZATION OF A SAMPLE OF THE SUBBAND
*                                     BEING PROCESSED
*                                     -2,3 or 4 BIT PCM CODING OF THAT QUANTIZED
*                                     SAMPLE
*
* INPUT: FAW + 1 kHz SAMPLES (8 STREAMS)
* OUTPUT: FAW + 1 kHz SAMPLES (8 STREAMS)
*-----*
* BIT ALLOCATION ACCORDING TO SSTRAT :
*
*
*          1  2  3  4  5  6  7  8
*-----*
*          VOICED      4  4  3  2  2  0  0  0
*          INTERMEDIATE 4  3  2  2  2  2  0  0
*          UNVOICED    2  3  3  3  2  2  0  0
*-----*
*
* AUTHOR : ROLAND KLEUTERS
* DATE : 29-5-1987
*****
*
PCMA8B NOP          ENTRY POINT
*
*   READ IN FROM TABLE THE NUMBER OF CODE BITS FOR THE SUBBAND IN PROCESSION
*   LACK BAV011
*   ADD PATH
*   ADD SSTRAT,3
*   TBLR NBIT          NBIT CONTAINS NUMBER OF CODE BITS
*
*   READ IN THE LOWEST AND HIGHEST POSSIBLE INPUT VALUE OF THE EXPONENT TABLE
*   LACK XX10
*   ADD PATH,1
*   TBLR X0          X0 CONTAINS LOWEST POSSIBLE INPUT OF EXP. TABLE
*   ADD ONE
*   TBLR X127       X127 CONTAINS HIGHEST POSSIBLE INPUT OF EXP. TABLE
*
*   CHECK IF NUMBER OF CODE BITS (NBIT) GREATER THAN ZERO
*   LAC NBIT
*   BGZ READIN      CODE?
*
*   NO CODING HAS TO TAKE PLACE (NBIT=0)
*   ZAC          ZERO "CODED" RESULT; I(n):=0
*   SACL CALC1
*   LACK COD1
*   ADD PATH
*   TBLW CALC1      I(n) SAVED IN CODED SUBBAND SAMPLES TABLE
*   LACK DN1        MINIMIZE STEPSIZE; d(n):=X0
*   ADD PATH
*   TBLW X0          d(n) SAVED IN TABLE
*   RET
*

```

```

*-----*
* DETERMINE 1/STEPSIZE i.e. 1/DELTA(n) BY BACKWARD STEPSIZE ADAPTION
*-----*
*
*.....*
*                               DETERMINE d(n)=LOG(DELTA(n))
*.....*
*
* READ IN THE PREVIOUS CODED RESULT I(n-1) FROM TABLE
READIN  LACK COD1
        ADD  PATH
        TBLR CALC1          CALC1 CONTAINS THE PREVIOUS CODED RESULT I(n-1)
*
* DETERMINE  $\mu([I(n-1)])=LOG(M([I(n-1)]))$  FROM TABLES
        LAC  NBIT
        SUB  ONE,1
        SACL CALC2
        LACK LOG22
        ADD  CALC2,4
        ADD  CALC1
        TBLR RES1          RES1 CONTAINS  $\mu([I(n-1)])=LOG(M([I(n-1)]))$ 
*
* DETERMINE  $(1-GAMMA)*DC$  FROM TABLE
        LACK OGD1
        ADD  PATH
        TBLR CALC1          CALC1 CONTAINS  $(1-GAMMA)*DC$ 
*
* DETERMINE  $GAMMA*d(n-1)$  FROM TABLE
        LACK DNI
        ADD  PATH
        TBLR RES0          RES0 CONTAINS  $d(n-1)$ 
        LT   GAMMA
        MPY  RES0
        PAC
        SACL CALC2,1       CALC2 CONTAINS  $GAMMA*d(n-1)$ 
*
* DETERMINE  $GAMMA*d(n-1)+\mu([I(n-1)])+(1-GAMMA)*DC$ 
        ZALH CALC2
        ADD  RES1,12
        ADD  CALC1,8
        SACL CALC1          CALC1 CONTAINS  $GAMMA*d(n-1)+\mu([I(n-1)])+$ 
*                                $(1-GAMMA)*DC$ , i.e.  $d(n)$ 
*
* CONTROL THE LIMITATION OF d(n) TO ITS RANGE AND SAVE d(n) IN TABLE
        LAC  CALC1
        SUB  X127
        B6Z SATDNH          d(n) GREATER THAN MAXIMUM POSSIBLE VALUE OF d(n)?
        LAC  X0
        SUB  CALC1
        B6Z SATDNL          d(n) LESS THAN MINIMUM POSSIBLE VALUE OF d(n)?
        B   SAVEDN          NO SATURATION NEEDED
SATDNH  LAC  X127          SATURATE d(n) TO MAXIMUM POSSIBLE VALUE
        SACL CALC1
        B   SAVEDN
SATDNL  LAC  X0           SATURATE d(n) TO MINIMUM POSSIBLE VALUE
        SACL CALC1

```

```

SAVEDN LACK DNI          SAVE d(n) IN TABLE
      ADD PATH
      TBLW CALC1
*.....
*          DETERMINE 1/DELTA(n)=EXP(-d(n))
*.....
*
*   ROUND d(n) TO ONE OF THE ENTRY POINTS OF THE EXPONENT TABLE
      ZALH CALC1
      SUBH X0
      SACH CALC1          CALC1 CONTAINS d(n)-X0
      LT CALC1
      MPY STEP
      PAC
      SACH CALC1          CALC1 CONTAINS (d(n)-X0)*STEP
      LAC CALC1,11
      SACH CALC1          CALC1 CONTAINS INTEGER[(d(n)-X0)*STEP]=ENTRY POINT
*
*   DETERMINE THE OUTPUT OF EXPONENT TABLE
      LACK CEXPO
      ADD CALC1
      TBLR CALC1          CALC1 CONTAINS OUTPUT OF EXPONENT TABLE=1/DELTA(n)
*
*-----
* QUANTIZE THE PCMA88 CODER INPUT AND CODE THE RESULT (DETERMINE I(n))
*-----
*
*   DETERMINE INPUT/DELTA(n) AND ROUND TO THE GREATEST INTEGER BELOW
      LT CALC1
      MPY INPCM
      PAC
      SACH CALC1          CALC1 CONTAINS INPUT/DELTA (Q18)
      LAC CALC1,14
      SACH CALC1          CALC1 CONTAINS INTEGER[INPUT/DELTA] IN
*                          TWO'S COMPLEMENT CODE, i.e. THE CODED RESULT I(n)
*
*   CONTROL THE LIMITATION OF I(n) TO THE RANGE ACCORDING TO THE NUMBER OF
*   CODE BITS AND SAVE I(n) IN TABLE OF CODED SUBBAND SAMPLES
      LACK SAT1
      ADD NBIT
      SUB ONE,1
      TBLR RES0          RES0 CONTAINS THE MAXIMUM POSSIBLE VALUE OF I(n)
      LAC CALC1
      SUB RES0
      B6Z SATINH          I(n) GREATER THAN MAXIMUM POSSIBLE VALUE OF I(n)?
      ZAC
      SUB RES0
      SUB ONE
      SACL RES1          RES1 CONTAINS THE MINIMUM POSSIBLE VALUE OF I(n)
      SUB CALC1
      B6Z SATINL          I(n) LESS THAN MINIMUM POSSIBLE VALUE OF I(n)?
      LACK COD1          SAVE I(n) WITHOUT SATURATION
      ADD PATH
      TBLW CALC1
      RET

```

```
SATINH LACK COD1      SATURATE I(n) TO MAXIMUM POSSIBLE VALUE AND SAVE
      ADD PATH
      TBLW RES0
      RET
SATINL LACK COD1      SATURATE I(n) TO MINIMUM POSSIBLE VALUE AND SAVE
      ADD PATH
      TBLW RES1
      RET
```

*

*

PCMAQB CODING DONE

=====

APPENDIX E
Subroutine MULTIPLEXING


```

*****
*                                     *
*                               MULTIPLEXING                               *
*                                     *
* THIS SUBROUTINE INCORPORATES:                                           *
*                                     *
*                               -MULTIPLEXING AND                          *
*                               -SENDING OF PCM CODED DATA                *
*                                     *
* INPUT: FAW + 1 kHz SAMPLES (8 STREAMS)                                  *
* OUTPUT: 16 KBPS BITSTREAM                                               *
*-----*
* SUBBAND PROCESSING : 1,8,4,5,2,7,3,6                                     *
*                                     *
* BIT ALLOCATION :                                                         *
*                                     *
*                               1  2  3  4  5  6  7  8                     *
*-----*-----*-----*-----*-----*-----*-----*-----*
*                               VOICED  4  4  3  2  2  0  0  0             *
*                               INTERMEDIATE 4  3  2  2  2  2  0  0       *
*                               UNVOICED  2  3  3  3  2  2  0  0         *
*                                     *
* FAW's :                                                                 *
*                               VOICED >76   INTERMEDIATE >96   UNVOICED >CC *
*-----*
* AUTHOR: ROLAND KLEUTERS                                               *
* DATE: 29-5-1987                                                       *
*****
*
*-----*
* MACRO
*-----*
*
*   MACRO TO SEND ONE BIT
SEND  $MACRO
      LAC DATA,1
      SACL DATA          SENDBIT PLACED IN FRONT OF BITSELECTOR
      LAC ONE,8
      AND DATA           SELECT SENDBIT
      SACL SIGOUT         SENDBIT IS BIT 8 OF SIGOUT
      OUT SIGOUT,3       SENDBIT SENDED VIA PIN 38 OF EXPANSION PORT
      $END
*
*-----*
* MACRO DEFINITION DONE; MAIN PART OF SUBROUTINE BEGINS
*-----*
*
MUX   NOP          ENTRY POINT
*
*   SELECT FAW OR ONE OF THE SUBBANDS FROM WHICH DATA HAS TO BE SENT
FAWSUB LACK 4
      SUB BAND
      B6Z FWSB6Z
      BZ NEXT4          SUBBAND 4 (1500-2000 Hz)
      ADD ONE
      BZ NEXT5          SUBBAND 5 (2000-2500 Hz)
      B NEXT6          SUBBAND 6 (2500-3000 Hz)

```

```
FWSB6Z SUB ONE
      BZ NEXT3          SUBBAND 3 (1000-1500 Hz)
      SUB ONE
      BZ NEXT2          SUBBAND 2 (500-1000 Hz)
      SUB ONE
      BZ NEXT1          SUBBAND 1 (0-500 Hz)
*
*
*-----*
* SEND FAW BITS
*-----*
*
NFAW  ZALS DEMO          NEW ACCESS CHECK
      BZ FSIG
*
*
*   INITIALIZATION FOR A NEW FAW ACCESS
      ZAC
      SACL DEMO          NO NEW ACCESS ANYMORE
      LACK 7
      SACL COUNT          8 FAW BITS TO SEND
*
*
*   LOAD FAW INTO SENDBUFFER (=DATA)
      ZALS SSTRAT
      BZ SVOICE
      XOR ONE
      BZ SINTER
SUNVOI LACK >CC          CODE FOR UNVOICED STRATEGY
      SACL DATA          PRESET UNVOICED STRATEGY
      B FSIG
SVOICE LACK >76          CODE FOR VOICED STRATEGY
      SACL DATA          PRESET VOICED STRATEGY
      B FSIG
SINTER LACK >96          CODE FOR INTERMEDIATE STRATEGY
      SACL DATA          PRESET INTERMEDIATE STRATEGY
*
*
*   SEND A FAW BIT
FSIG  SEND
      LAR 0,COUNT
      BANZ CONT
*
*
*   ALL FAW BITS SENT,PREPARE NEXT SUBBAND
      LACK >FF
      SACL DEMO          NEW ACCESS MARK
      LACK 1
      SACL BAND          NEXT SUBBAND TO SEND (0-500 Hz)
      RET
*
*
*   NOT ALL FAW BITS SEND,PREPARE NEXT BIT
CONT  SAR 0,COUNT
      RET
*
*-----*
* SEND SUBBAND BITS
*-----*
*
```

```
*.....
*           A BIT OF SUBBAND 1 (0-500 Hz) HAS TO BE SENT
*.....
*
NEXT1  ZALS DEMO           NEW ACCESS CHECK
      BZ  LOOP1
*
*  INITIALIZATION FOR A NEW SUBBAND 1 ACCESS
      ZAC
      SACL DEMO           NO NEW ACCESS ANYMORE
      LACK COD1
      TBLR DATA         LOAD SUBBAND 1 DATA INTO SEND BUFFER
      LACK 2
      XOR SSTRAT         UNVOICED?
      BNZ VOINI
*
*  FURTHER INITIALIZATION FOR UNVOICED STRATEGY
      LACK 1
      SACL COUNT         2 BITS TO SEND
      LAC DATA,6       ADJUST DATA ACCORDING TO BITSELECTOR
      SACL DATA         DATA=0000 0000 D000 0000
      B  LOOP1
*
*  FURTHER INITIALIZATION FOR VOICED AND INTERMEDIATE STRATEGY
VOINI  LACK 3
      SACL COUNT         4 BITS TO SEND
      LAC DATA,4       ADJUST DATA ACCORDING TO BITSELECTOR
      SACL DATA         DATA=0000 0000 DDDD 0000
*
*  SEND A SUBBAND 1 BIT
LOOP1  SEND
      LAR 0,COUNT
      BANZ CONT1
*
*  ALL SUBBAND 1 BITS SENT,PREPARE NEXT SUBBAND
      LACK >FF
      SACL DEMO           NEW ACCESS MARK
      LACK 4
      SACL BAND           NEXT SUBBAND TO SEND (1500-2000 Hz)
      RET
*
*  NOT ALL SUBBAND 1 BITS SENT,PREPARE NEXT BIT
CONT1  SAR 0,COUNT
      RET
*
*.....
*           A BIT OF SUBBAND 2 (500-1000 Hz) HAS TO BE SENT
*.....
*
NEXT2  ZALS DEMO           NEW ACCESS CHECK
      BZ  LOOP2
*
```

```
*   INITIALIZATION FOR A NEW SUBBAND 2 ACCESS
    ZAC
    SACL DEMO           NO NEW ACCESS ANYMORE
    LACK COD2
    TBLR DATA         LOAD SUBBAND 2 DATA INTO SEND BUFFER
    LAC SSTRAT         VOICED?
    BNZ INUN2
*
*   FURTHER INITIALIZATION FOR VOICED STRATEGY
    LACK 3
    SACL COUNT         4 BITS TO SEND
    LAC DATA,4        ADJUST DATA ACCORDING TO BITSELECTOR
    SACL DATA         DATA=0000 0000 DDDD 0000
    B LOOP2
*
*   FURTHER INITIALIZATION FOR INTERMEDIATE AND UNVOICED STRATEGY
INUN2 LACK 2
    SACL COUNT         3 BITS TO SEND
    LAC DATA,5        ADJUST DATA ACCORDING TO BITSELECTOR
    SACL DATA         DATA=0000 0000 DDDD 0000
*
*   SEND A SUBBAND 2 BIT
LOOP2 SEND
    LAR 0,COUNT
    BANZ CONT2
*
*   ALL SUBBAND 2 BITS SENT,PREPARE NEXT SUBBAND
    LACK >FF
    SACL DEMO           NEW ACCESS MARK
    LACK 3
    SACL BAND           NEXT SUBBAND TO SEND (1000-1500 Hz)
    RET
*
*   NOT ALL SUBBAND 2 BITS SENT,PREPARE NEXT BIT
CONT2 SAR 0,COUNT
    RET
*
* .....
*           A BIT OF SUBBAND 3 (1000-1500) HAS TO BE SENT
* .....
*
NEXT3 ZALS DEMO         NEW ACCESS CHECK
    BZ LOOP3
*
*   INITIALIZATION FOR A NEW SUBBAND 3 ACCESS
    ZAC
    SACL DEMO           NO NEW ACCESS ANYMORE
    LACK COD3
    TBLR DATA         LOAD SUBBAND 3 DATA INTO SEND BUFFER
    LACK 1
    XDR SSTRAT         INTERMEDIATE?
    BNZ VOUN3
*
```

```
* FURTHER INITIALIZATION FOR INTERMEDIATE STRATEGY
  LACK 1
  SACL COUNT      2 BITS TO SEND
  LAC DATA,6    ADJUST DATA ACCORDING TO BITSELECTOR
  SACL DATA      DATA=0000 0000 D000 0000
  B LOOP3
*
* FURTHER INITIALIZATION FOR VOICED AND UNVOICED STRATEGY
VOUN3 LACK 2
      SACL COUNT      3 BITS TO SEND
      LAC DATA,5    ADJUST DATA ACCORDING TO BITSELECTOR
      SACL DATA      DATA=0000 0000 D000 0000
*
* SEND A SUBBAND 3 BIT
LOOP3 SEND
      LAR 0,COUNT
      BANZ CONT3
*
* ALL SUBBAND 3 BITS SENT AND PERHAPS ALSO END OF 15 BIT DATABLOCK
* CHECK FOR END OF FRAME
      LACK >FF
      SACL DEMO      NEW ACCESS MARK
      LACK 7
      XOR FRAME
      BZ ENDFRM
*
* NOT END OF FRAME,PREPARE NEXT SUBBAND
      LAC SSTRAT      VOICED?
      BZ MENDFM
      LACK 6
      SACL BAND      NEXT SUBBAND TO SEND (2500-3000 Hz)
      RET
*
* NOT ALL SUBBAND 3 BITS SENT,PREPARE NEXT BIT
CONT3 SAR 0,COUNT
      RET
*
*.....
*                A BIT OF SUBBAND 4 (1500-2000 Hz) HAS TO BE SENT
*.....
*
NEXT4 ZALS DEMO      NEW ACCESS CHECK
      BZ LOOP4
*
* INITIALIZATION FOR A NEW SUBBAND 4 ACCESS
      ZAC
      SACL DEMO      NO NEW ACCESS ANYMORE
      LACK COD4
      TBLR DATA    LOAD SUBBAND 4 DATA INTO SEND BUFFER
      LACK 2
      XOR SSTRAT    UNVOICED?
      BNZ VOIN4
*
```

```
* FURTHER INITIALIZATION FOR UNVOICED STRATEGY
  LACK 2
  SACL COUNT      3 BITS TO SEND
  LAC DATA,5    ADJUST DATA ACCORDING TO BITSELECTOR
  SACL DATA      DATA=0000 0000 D000 0000
  B LOOP4
*
* FURTHER INITIALIZATION FOR VOICED AND INTERMEDIATE STRATEGY
VOIN4 LACK 1
      SACL COUNT      2 BITS TO SEND
      LAC DATA,6    ADJUST DATA ACCORDING TO BITSELECTOR
      SACL DATA      DATA=0000 0000 D000 0000
*
* SEND A SUBBAND 4 BIT
LOOP4 SEND
      LAR 0,COUNT
      BANZ CONT4
*
* ALL SUBBAND 4 BITS SENT PREPARE NEXT SUBBAND
      LACK >FF
      SACL DEMO        NEW ACCESS MARK
      LACK 5
      SACL BAND        NEXT SUBBAND TO SEND (2000-2500 Hz)
      RET
*
* NOT ALL SUBBAND 4 BITS SENT,PREPARE NEXT BIT
CONT4 SAR 0,COUNT
      RET
*
*.....
*          A BIT OF SUBBAND 5 (2000-2500) HAS TO BE SENT
*.....
*
NEXT5 ZALS DEMO        NEW ACCESS CHECK
      BZ LOOP5
*
* INITIALIZATION FOR A NEW SUBBAND 5 ACCESS
      ZAC
      SACL DEMO        NO NEW ACCESS ANYMORE
      LACK C005
      TBLR DATA      LOAD SUBBAND 5 DATA INTO SEND BUFFER
      LACK 1
      SACL COUNT      2 BITS TO SEND FOR ALL STRATEGIES
      LAC DATA,6    ADJUST DATA ACCORDING TO BITSELECTOR
      SACL DATA      DATA=0000 0000 D000 0000
*
* SEND A SUBBAND 5 BIT
LOOP5 SEND
      LAR 0,COUNT
      BANZ CONT5
*
* ALL SUBBAND 5 BITS SENT,PREPARE NEXT SUBBAND
      LACK >FF
      SACL DEMO        NEW ACCESS MARK
      LACK 2
      SACL BAND        NEXT SUBBAND TO SEND (500-1000 Hz)
      RET
*
```

```
* NOT ALL SUBBAND 5 BITS SENT,PREPARE NEXT BIT
CONT5 SAR 0,COUNT
RET
*
*.....
* A BIT OF SUBBAND 6 (2500-3000 Hz) HAS TO BE SENT
* ONLY ACCESSED BY UNVOICED AND INTERMEDIATE STRATEGY
*.....
*
NEXT6 ZALS DEMO NEW ACCESS CHECK
BZ LOOP6
*
* INITIALIZATION FOR A NEW SUBBAND 6 ACCESS
ZAC
SACL DEMO NO NEW ACCESS ANYMORE
LACK COD6
TBLR DATA LOAD SUBBAND 6 DATA INTO SEND BUFFER
LACK 1
SACL COUNT 2 BITS TO SEND FOR ALL POSSIBLE STRATEGIES
LAC DATA,6 ADJUST DATA ACCORDING TO BITSELECTOR
SACL DATA DATA=0000 0000 DD00 0000
*
* SEND A SUBBAND 6 BIT
LOOP6 SEND
LAR 0,COUNT
BANZ CONT6
*
* ALL SUBBAND 6 BITS SENT AND ALSO END OF 15 BIT DATABLOCK
* CHECK FOR END OF FRAME
LACK >FF
SACL DEMO NEW ACCESS MARK
LACK 7
XOR FRAME
BNZ NENDFM
*
* END OF FRAME,PREPARE NEXT FRAME
ENDFM ZAC
SACL BAND NEXT TO SEND=FAW
RET
*
* NOT END OF FRAME,PREPARE NEXT DATABLOCK
NENDFM LACK 1
SACL BAND NEXT SUBBAND TO SEND (0-500 Hz)
RET
*
* NOT ALL SUBBAND 6 BITS SENT,PREPARE NEXT BIT
CONT6 SAR 0,COUNT
RET
*
* MULTIPLEXING DONE
*=====
```

APPENDIX F

Main program SBC receiver part


```
*****
*                               DECODER PROGRAM                               *
*                               *                                           *
* THIS MAIN PROGRAM USES THE SUBROUTINES:                                  *
*                               -DEMULPLEXING                               *
*                               -PCMAQR DECODING                           *
*                               -INVQMF FILTERING                           *
* TO REALIZE A 16 KBIT PER SECOND SUBBAND DECODER.                         *
* SYSTEM CLOCK RATE: 16 kHz                                              *
*-----*
* REPORTS FOR DETAILS:                                                    *
* PT REPORT BY P.CHITAMU                                                  *
* FINAL REPORT BY J.BOLT.                                                 *
* FINAL REPORT BY R.KLEUTERS                                              *
*-----*
* AUTHOR: ROLAND KLEUTERS                                                *
* DATE: 29-5-1987                                                         *
*****
*
*-----*
* DATA MEMORY INITIALIZATION                                           *
*                               default page=0                             *
*-----*
*.....*
*                               GENERAL VARIABLES DATA MEMORY PAGE 1     *
*.....*
*
CLOCK EQU 0          CLOCK RATE (16 kHz)
MODE EQU 1          MODE FOR ANALOG INTERFACE BOARD
AR00 EQU 2          STORAGE PLACE FOR AUXILIARY REGISTER 0
AR01 EQU 3          STORAGE PLACE FOR AUXILIARY REGISTER 1
ACH EQU 4           STORAGE PLACE FOR HIGH PART OF ACCUMULATOR
ACL EQU 5           STORAGE PLACE FOR LOW PART OF ACCUMULATOR
TREG EQU 6          STORAGE PLACE FOR T REGISTER
STATU EQU 7         STORAGE PLACE FOR STATUS WORD OF TMS32010
*                               from now on default page
*.....*
*                               GENERAL VARIABLES                           *
*.....*
*
MSTAT EQU 0         MASK FOR STATE MOD 16
SWITCH EQU 1        >FF=SEARCH MODE,ELSE=RECEIVE MODE; 0=DECODER ON(Q0)
FRAME EQU 2         COUNTER FOR FRAMELENGTH (Q0)
ONE EQU 3           CHECK BIT (Q0)
STRAT EQU 4         CONTAINS 0, 1, or 2: VOICED, INTERM. or UNVD. (Q0)
SSTRAT EQU 5        BUFFERING (Q0)
STATE EQU 6         TREE POINTER (Q0)
RES0 EQU 7          INTERMEDIATE CALCULATION RESULTS FOR MAIN ROUTINE
RES1 EQU 8          INTERMEDIATE CALCULATION RESULTS FOR MAIN ROUTINE
PATH EQU 9          CONTAINS STATE/2=CHANNEL TO PROCESS (Q0)
OUTP EQU 10         8 kHz OUTPUT SAMPLE OF TMS32010 (Q15)
BOUTP1 EQU 11       BUFFERED 8 kHz OUTPUT SAMPLE OF TMS32010 (Q15)
BOUTP2 EQU 12       16 kHz INTERPOLATION OF TMS32010 OUTPUT (Q15)
*

```

*.....
* FIRST INVQMF STAGE VARIABLES
*.....

* BUFFER VARIABLES OF LOW1 (0-2 kHz) AND HIGH1 (2-4 kHz) FOR $\omega=2:4$ RESP. Q16

A1 EQU 13
B1 EQU 14
C1 EQU 15

* INPUT TO (FILTER PART OF) FIRST INVQMF STAGE FORMAT=Q15
INPF EQU 16

* DELAY VARIABLES OF LOW1 (0-2 kHz) FORMAT=Q15

XF2 EQU 17 AFTER FILTERING INPUT STORED IN XF2, SO
XF3 EQU 18 XF1 NOT NEEDED
XF4 EQU 19
XF5 EQU 20
XF6 EQU 21
XF7 EQU 22
XF8 EQU 23
XF9 EQU 24
XF10 EQU 25
XF11 EQU 26
XF12 EQU 27
XF13 EQU 28
XF14 EQU 29
XF15 EQU 30
XF16 EQU 31

* DELAY VARIABLES OF HIGH1 (2-4 kHz) FORMAT=Q15

XF18 EQU 32 AFTER FILTERING INPUT STORED IN XF18, SO
XF19 EQU 33 XF17 NOT NEEDED
XF20 EQU 34
XF21 EQU 35
XF22 EQU 36
XF23 EQU 37
XF24 EQU 38
XF25 EQU 39
XF26 EQU 40
XF27 EQU 41
XF28 EQU 42
XF29 EQU 43
XF30 EQU 44
XF31 EQU 45
XF32 EQU 46

*.....
* SECOND INVQMF STAGE VARIABLES
*.....

* BUFFER VARIABLES OF LOW2 (0-1 kHz) AND HIGH2 (1-2 kHz) FORMAT=Q14

A2 EQU 47
B2 EQU 48
C2 EQU 49

*

```
* BUFFER VARIABLES OF LOW3 (3-4 kHz) AND HIGH3 (2-3 kHz)          FORMAT=Q17
A3   EQU 50
B3   EQU 51
C3   EQU 52
*
* INPUT TO (FILTER PART OF) SECOND INVQMF STAGE
INPS EQU 53
*
* DELAY VARIABLES OF LOW2,HIGH2,LOW3 AND HIGH3 (TIME SHARING)
XS2  EQU 54
XS3  EQU 55
XS4  EQU 56
XS5  EQU 57
XS6  EQU 58
XS7  EQU 59
XS8  EQU 60
*
*.....
*                               THIRD INVQMF STAGE VARIABLES
*.....
*
* BUFFER VARIABLES OF LOW4 (0-500 Hz) AND HIGH4 (500-1000 Hz)    FORMAT=Q14
A4   EQU 61
B4   EQU 62
C4   EQU 63
*
* BUFFER VARIABLES OF LOW5 (1500-2000 Hz) AND HIGH5 (1000-1500 Hz) FOR.=Q15
A5   EQU 64
B5   EQU 65
C5   EQU 66
*
* BUFFER VARIABLES OF LOW7 (2000-2500 Hz) AND HIGH7 (2500-3000 Hz) FOR.=Q18
A7   EQU 67
B7   EQU 68
C7   EQU 69
*
* INPUT TO (FILTER PART OF) THIRD INVQMF STAGE
INPT EQU 70
*
* DELAY VARIABLES OF LOW4,HIGH4,LOW5,HIGH5,LOW7 AND HIGH7 (TIME SHARING)
XT2  EQU 71
XT3  EQU 72
XT4  EQU 73
XT5  EQU 74
XT6  EQU 75
*
*.....
*                               FIRST INVQMF STAGE COEFFICIENTS          FORMAT=Q16
*.....
*
CF0  EQU 76
CF1  EQU 77
CF2  EQU 78
CF3  EQU 79
CF4  EQU 80
CF5  EQU 81
CF6  EQU 82
CF7  EQU 83
```

CF8 EQU 84
CF9 EQU 85
CF10 EQU 86
CF11 EQU 87
CF12 EQU 88
CF13 EQU 89
CF14 EQU 90
CF15 EQU 91

*

*.....
* SECOND INVQMF STAGE COEFFICIENTS FORMAT=Q16
*.....

CS0 EQU 92
CS1 EQU 93
CS2 EQU 94
CS3 EQU 95
CS4 EQU 96
CS5 EQU 97
CS6 EQU 98
CS7 EQU 99

*

*.....
* THIRD INVQMF STAGE COEFFICIENTS FORMAT=Q16
*.....

CT0 EQU 100
CT1 EQU 101
CT2 EQU 102
CT3 EQU 103
CT4 EQU 104
CT5 EQU 105

*

*.....
* PCMAQB VARIABLES
*.....

OUTPCM EQU 106 OUTPUT OF PCMAQB DECODER
CALC1 EQU 107 GENERAL VARIABLE TO CALCULATE WITH
CALC2 EQU 108 GENERAL VARIABLE TO CALCULATE WITH
NBIT EQU 109 CONTAINS NR. OF BITS TO QUANTIZE TO (Q0)
X0 EQU 110 LOWEST POSSIBLE INPUT VALUE OF EXPONENT TABLE (Q14)
X127 EQU 111 HIGHEST POSSIBLE INPUT VALUE OF EXPONENT TABLE(Q14)
STEP EQU 112 1/(DISTANCE BETWEEN INPUTS OF EXPONENT TABLE) (Q7)
GAMMA EQU 113 GAMMA VALUE OF FORMULA (0.99) (Q15)
V3FFF EQU 114 >3FFF; HELPCONSTANT

*

*.....
* DEMULTIPLEXER VARIABLES
*.....

POINT EQU 115 FLAG: >FF=SEARCH FAW, 0=SKIP 128 BITS
ERROR EQU 116 CONTAINS NUMBER OF ERROR ALLOWANCES (Q0)
COUNT EQU 117 MULTIPURPOSE COUNTER (Q0)
BAND EQU 118 FLAG, CURRENT SUBBAND PROCESSING (Q0)
DATA EQU 119 STORES INPUT BITS VIA SIGIN
SIGIN EQU 120 VARIABLE TO READ IN SERIAL DATA FROM OUTSIDE
DEND EQU 121 NEW (FAW or SUBBAND) ACCESS MARK

```
*
*-----*
* PROGRAM MEMORY INITIALIZATION
*-----*
*
*      AORG 0
*      B   START          BRANCH TO MAIN ROUTINE
*      B   INTRO         BRANCH TO INTERRUPT SERVICE ROUTINE
*      AORG 8             OTHERWISE ERRORS WITH TBLW'S
*
*.....*
*                          GENERAL CONSTANTS
*.....*
*
*      CCLOCK DATA 312      CLOCK RATE (16 kHz)
*      MMODE DATA 7        MODE FOR ANALOG INTERFACE BOARD
*      MMSTAT DATA >000F   MASK FOR STATE MOD 16
*
*.....*
*                          SECOND INVQMF STAGE VARIABLES
*.....*
*
*      DELAY VARIABLES OF LOW2 (0-1 kHz)                                FORMAT=Q14
*      XXS2 DATA 0          AFTER FILTERING INPUT STORED IN XXS2, S0
*      XXS3 DATA 0          XXS1 NOT NEEDED
*      XXS4 DATA 0
*      XXS5 DATA 0
*      XXS6 DATA 0
*      XXS7 DATA 0
*      XXS8 DATA 0
*
*      DELAY VARIABLES OF HIGH2 (1-2 kHz)                                FORMAT=Q14
*      XXS10 DATA 0         AFTER FILTERING INPUT STORED IN XXS10, S0
*      XXS11 DATA 0         XXS9 NOT NEEDED
*      XXS12 DATA 0
*      XXS13 DATA 0
*      XXS14 DATA 0
*      XXS15 DATA 0
*      XXS16 DATA 0
*
*      DELAY VARIABLES OF LOW3 (3-4 kHz)                                FORMAT=Q16
*      XXS18 DATA 0         AFTER FILTERING INPUT STORED IN XXS18, S0
*      XXS19 DATA 0         XXS17 NOT NEEDED
*      XXS20 DATA 0
*      XXS21 DATA 0
*      XXS22 DATA 0
*      XXS23 DATA 0
*      XXS24 DATA 0
*
*      DELAY VARIABLES OF HIGH3 (2-3 kHz)                                FORMAT=Q16
*      XXS26 DATA 0         AFTER FILTERING INPUT STORED IN XXS26, S0
*      XXS27 DATA 0         XXS25 NOT NEEDED
*      XXS28 DATA 0
*      XXS29 DATA 0
*      XXS30 DATA 0
*      XXS31 DATA 0
*      XXS32 DATA 0
*
```

*.....
* THIRD INVQMF STAGE VARIABLES
*.....

* DELAY VARIABLES OF LOW4 (0-500 Hz) FORMAT=Q14

XXT2 DATA 0 AFTER FILTERING INPUT STORED IN XXT2, SO
XXT3 DATA 0 XXT1 NOT NEEDED
XXT4 DATA 0
XXT5 DATA 0
XXT6 DATA 0

* DELAY VARIABLES OF HIGH4 (500-1000 Hz) FORMAT=Q14

XXT8 DATA 0 AFTER FILTERING INPUT STORED IN XXT8, SO
XXT9 DATA 0 XXT7 NOT NEEDED
XXT10 DATA 0
XXT11 DATA 0
XXT12 DATA 0

* DELAY VARIABLES OF LOW5 (1500-2000 Hz) FORMAT=Q14

XXT14 DATA 0 AFTER FILTERING INPUT STORED IN XXT14, SO
XXT15 DATA 0 XXT13 NOT NEEDED
XXT16 DATA 0
XXT17 DATA 0
XXT18 DATA 0

* DELAY VARIABLES OF HIGH5 (1000-1500 Hz) FORMAT=Q14

XXT20 DATA 0 AFTER FILTERING INPUT STORED IN XXT20, SO
XXT21 DATA 0 XXT19 NOT NEEDED
XXT22 DATA 0
XXT23 DATA 0
XXT24 DATA 0

* DELAY VARIABLES OF LOW7 (2000-2500 Hz) FORMAT=Q17

XXT26 DATA 0 AFTER FILTERING INPUT STORED IN XXT26, SO
XXT27 DATA 0 XXT25 NOT NEEDED
XXT28 DATA 0
XXT29 DATA 0
XXT30 DATA 0

* DELAY VARIABLES OF HIGH7 (2500-3000 Hz) FORMAT=Q17

XXT32 DATA 0 AFTER FILTERING INPUT STORED IN XXT32, SO
XXT33 DATA 0 XXT31 NOT NEEDED
XXT34 DATA 0
XXT35 DATA 0
XXT36 DATA 0

*.....
* FIRST INVQMF STAGE COEFFICIENTS FORMAT=Q16
*.....

CCF0 DATA 45
CCF1 DATA -92
CCF2 DATA -83
CCF3 DATA 277
CCF4 DATA 93
CCF5 DATA -620
CCF6 DATA -9
CCF7 DATA 1178

CCF8 DATA -274
CCF9 DATA -2047
CCF10 DATA 955
CCF11 DATA 3470
CCF12 DATA -2579
CCF13 DATA -6541
CCF14 DATA 8425
CCF15 DATA 30566

*

*.....
* SECOND INVQMF STAGE COEFFICIENTS FORMAT=Q16
*.....

*
CCS0 DATA 69
CCS1 DATA -331
CCS2 DATA -170
CCS3 DATA 1812
CCS4 DATA -633
CCS5 DATA -5924
CCS6 DATA 6409
CCS7 DATA 31525

*

*.....
* THIRD INVQMF STAGE COEFFICIENTS FORMAT=Q16
*.....

*
CCT0 DATA -250
CCT1 DATA 1236
CCT2 DATA -178
CCT3 DATA -5551
CCT4 DATA 5798
CCT5 DATA 31745

*

*.....
* PCMAQB'S TABLES
*.....

* (1-GAMMA)*DC FOR 7 SUBBANDS FORMAT=Q22
O6D1 DATA -10486 SUBBAND 1
DUM1 DATA 0
O6D4 DATA -13642 SUBBAND 4
O6D5 DATA -23112 SUBBAND 5
O6D2 DATA -10486 SUBBAND 2
DUM2 DATA 0
O6D3 DATA -13642 SUBBAND 3
O6D6 DATA -23112 SUBBAND 6

*

* D(N) FOR 7 SUBBANDS FORMAT=Q14
DN1 DATA -13563 SUBBAND 1
DUM3 DATA 0
DN4 DATA -14796 SUBBAND 4
DN5 DATA -18495 SUBBAND 5
DN2 DATA -13563 SUBBAND 2
DUM4 DATA 0
DN3 DATA -14796 SUBBAND 3
DN6 DATA -18495 SUBBAND 6

*

```
* BIT ALLOCATION FOR VOICED STRATEGY (0)                                FORMAT=Q0
BAVO11 DATA 4 SUBBAND 1
DUM5 DATA 0
BAVO14 DATA 2 SUBBAND 4
BAVO15 DATA 2 SUBBAND 5
BAVO12 DATA 4 SUBBAND 2
DUM6 DATA 0
BAVO13 DATA 3 SUBBAND 3
BAVO16 DATA 0 SUBBAND 6
*
* BIT ALLOCATION FOR INTERMEDIATE STRATEGY (1)                        FORMAT=Q0
BAINT1 DATA 4 SUBBAND 1
DUM7 DATA 0
BAINT4 DATA 2 SUBBAND 4
BAINT5 DATA 2 SUBBAND 5
BAINT2 DATA 3 SUBBAND 2
DUM8 DATA 0
BAINT3 DATA 2 SUBBAND 3
BAINT6 DATA 2 SUBBAND 6
*
* BIT ALLOCATION FOR UNVOICED STRATEGY (2)                            FORMAT=Q0
BAUNV1 DATA 2 SUBBAND 1
DUM9 DATA 0
BAUNV4 DATA 3 SUBBAND 4
BAUNV5 DATA 2 SUBBAND 5
BAUNV2 DATA 3 SUBBAND 2
DUM10 DATA 0
BAUNV3 DATA 3 SUBBAND 3
BAUNV6 DATA 2 SUBBAND 6
*
* STORAGE OF CODED SUBBAND SAMPLES (INPUT PCMAQB DECODER)           FORMAT=Q0
COD1 DATA 0 SUBBAND 1
DUM11 DATA 0
COD4 DATA 0 SUBBAND 4
COD5 DATA 0 SUBBAND 5
COD2 DATA 0 SUBBAND 2
DUM12 DATA 0
COD3 DATA 0 SUBBAND 3
COD6 DATA 0 SUBBAND 6
*
* STORAGE OF PREVIOUS CODED SUBBAND SAMPLES (STEP SIZE ADAPTION)    FORMAT=Q0
PCOD1 DATA 0 SUBBAND 1
DUM13 DATA 0
PCOD4 DATA 0 SUBBAND 4
PCOD5 DATA 0 SUBBAND 5
PCOD2 DATA 0 SUBBAND 2
DUM14 DATA 0
PCOD3 DATA 0 SUBBAND 3
PCOD6 DATA 0 SUBBAND 6
*
* QUANTIZER CONSTANTS
SSTEP DATA 19637 1/(DISTANCE BETWEEN INPUTS OF EXPONENT TABLE) (Q7)
66AMMA DATA 32440 GAMMA=0.99 (Q15)
VV3FFF DATA >3FFF >3FFF; HELPCONSTANT
*
```


* LOGTABLE 2 BIT CASE

FORMAT=Q18

LOG20 DATA 18268
LOG21 DATA -4626
LOG22 DATA -4626
LOG23 DATA 18268
DUM15 DATA 0
DUM16 DATA 0
DUM17 DATA 0
DUM18 DATA 0
DUM19 DATA 0
DUM20 DATA 0
DUM21 DATA 0
DUM22 DATA 0
DUM23 DATA 0
DUM24 DATA 0

*

* LOGTABLE 3 BIT CASE

FORMAT=Q18

LOG30 DATA 11540
LOG31 DATA 0
LOG32 DATA 0
LOG33 DATA -4626
LOG34 DATA -4626
LOG35 DATA 0
LOG36 DATA 0
LOG37 DATA 11540
DUM25 DATA 0
DUM26 DATA 0
DUM27 DATA 0
DUM28 DATA 0

*

* LOGTABLE 4 BIT CASE

FORMAT=Q18

LOG40 DATA 24918
LOG41 DATA 19728
LOG42 DATA 13377
LOG43 DATA 5189
LOG44 DATA -2999
LOG45 DATA -2999
LOG46 DATA -2999
LOG47 DATA -2999
LOG48 DATA -2999
LOG49 DATA -2999
LOG410 DATA -2999
LOG411 DATA -2999
LOG412 DATA 5189
LOG413 DATA 13377
LOG414 DATA 19728
LOG415 DATA 24918

*

* LOWEST AND HIGHEST POSSIBLE INPUTS OF EXPONENT TABLE

FORMAT=Q14

XX10 DATA -13563 SUBBAND 1
XX1127 DATA 0 SUBBAND 1
DUM29 DATA 0
DUM30 DATA 0
XX40 DATA -14796 SUBBAND 4
XX4127 DATA -1233 SUBBAND 4
XX50 DATA -18495 SUBBAND 5
XX5127 DATA -4932 SUBBAND 5
XX20 DATA -13563 SUBBAND 2

XX2127	DATA 0	SUBBAND 2
DUM31	DATA 0	
DUM32	DATA 0	
XX30	DATA -14796	SUBBAND 3
XX3127	DATA -1233	SUBBAND 3
XX60	DATA -18495	SUBBAND 6
XX6127	DATA -4932	SUBBAND 6

*
* EXPONENT TABLE OF DECODER
* FORMATS ARE : SUBBAND 1 Q15
* SUBBAND 2 Q15
* SUBBAND 3 Q16
* SUBBAND 4 Q16
* SUBBAND 5 Q19
* SUBBAND 6 Q19
*

CXPO DATA 16
DATA 17
DATA 18
DATA 19
DATA 20
DATA 22
DATA 23
DATA 24
DATA 26
DATA 27
DATA 29
DATA 31
DATA 33
DATA 35
DATA 37
DATA 39
DATA 42
DATA 44
DATA 47
DATA 50
DATA 53
DATA 56
DATA 60
DATA 64
DATA 68
DATA 72
DATA 76
DATA 81
DATA 86
DATA 91
DATA 97
DATA 103
DATA 109
DATA 116
DATA 123
DATA 131
DATA 139
DATA 148
DATA 157
DATA 166
DATA 177
DATA 188

DATA 199
DATA 211
DATA 225
DATA 238
DATA 253
DATA 269
DATA 286
DATA 303
DATA 322
DATA 342
DATA 363
DATA 385
DATA 409
DATA 435
DATA 462
DATA 490
DATA 520
DATA 553
DATA 587
DATA 623
DATA 662
DATA 703
DATA 746
DATA 792
DATA 841
DATA 893
DATA 949
DATA 1007
DATA 1070
DATA 1136
DATA 1206
DATA 1281
DATA 1360
DATA 1444
DATA 1534
DATA 1628
DATA 1729
DATA 1836
DATA 1950
DATA 2070
DATA 2199
DATA 2335
DATA 2479
DATA 2632
DATA 2795
DATA 2968
DATA 3152
DATA 3347
DATA 3554
DATA 3774
DATA 4008
DATA 4256
DATA 4519
DATA 4798
DATA 5095
DATA 5411
DATA 5745
DATA 6101

DATA 6478
DATA 6879
DATA 7305
DATA 7757
DATA 8237
DATA 8746
DATA 9288
DATA 9862
DATA 10473
DATA 11121
DATA 11809
DATA 12539
DATA 13315
DATA 14139
DATA 15014
DATA 15943
DATA 16929
DATA 17977
DATA 19089
DATA 20270
DATA 21525
DATA 22856
DATA 24271
DATA 25772
DATA 27367
DATA 29060
DATA 30859
DATA 32767

+

+ MAIN ROUTINE

+

+

.....
+ INITIALIZATIONS
.....

+

START DINT

+

+ INITIALIZE STATUS BITS

LARP 0
LDPK 0
SOVM

+

+ INITIALIZE GENERAL VARIABLES OF DATA MEMORY PAGE 1

LDPK 1
LACK CCLOCK
TBLR CLOCK CLOCK RATE (16 kHz)
LACK MMODE
TBLR MODE MODE FOR ANALOG INTERFACE BOARD
LDPK 0

+

+ INITIALIZE GENERAL VARIABLES OF DATA MEMORY PAGE 0

LACK MMSTAT
TBLR MSTAT MASK FOR STATE MOD 16
LACK >FF
SACL SWITCH DEMULTIPLEXER IN SEARCH MODE; DECODER OFF
ZAC

```
SACL FRAME          STATUS FOR DECODER
LACK 1
SACL ONE            '1' FOR LOGICAL MANIPULATIONS
SACL STRAT          BY DEFAULT LOAD INTERMEDIATE STRATEGY
SACL SSTRAT        ALSO INTERMEDIATE STRATEGY FOR BUFFER
ZAC
SACL STATE          INIT TREE POINTER
SACL OUTP           ZERO 8 kHz OUTPUT SAMPLE OF TMS32010
SACL BOUTP1        ZERO BUFFERED 8 kHz OUTPUT SAMPLE OF TMS32010
SACL BOUTP2        ZERO 16 kHz INTERPOLATED OUTPUT OF TMS32010
*
* INITIALIZE PCMAQB VARIABLES
  LACK SSTEP
  TBLR STEP          1/(DISTANCE BETWEEN INPUTS OF EXPONENT TABLE)
  LACK 6GAMMA
  TBLR 6GAMMA        GAMMA VALUE OF FORMULA (0.99)
  LACK VV3FFF
  TBLR V3FFF         >3FFF; HELPCONSTANT
*
* INITIALIZE DEMULTIPLEXER VARIABLES
  LACK >FF
  SACL POINT         SEARCH FIRST FAW
*
* LOAD INVQMF COEFFICIENTS INTO DATA RAM
  LARK 1,CTS
  LARK 0,29
  LACK CCTS
LOAD  LARP 1
      TBLR *-,0
      SUB ONE
      BANZ LOAD
*
* INITIALIZE ANALOG INTERFACE BOARD (AIB)
  LDPK 1
  OUT MODE,0         PROGRAM AIB
  OUT CLOCK,1        SET CLOCK RATE OF AIB
  LDPK 0
  OUT OUTP,2         CLEAR AD1S (PROTECTION AGAINST RESET INTERRUPTS)
  OUT OUTP,3         CLEAR AD2S (PROTECTION AGAINST RESET INTERRUPTS)
*
* INITIALIZATION AND VARIABLE LOADING DONE, NOW BEGIN
  EINT
*
*.....
*                                AORTA
*.....
*
* ODD CYCLE OF 16 kHz CLOCK MUST JUST HAVE PASSED
WAITEV LAC STATE
      AND ONE
      BNZ WAITEV
* EVEN CYCLE OF 16 kHz CLOCK HAS PASSED (BIT 0 OF STATE=0)
WAIT  LAC STATE
      AND ONE
      BZ WAIT
* ODD CYCLE OF 16 kHz CLOCK HAS JUST PASSED (BIT 0 OF STATE=1)
*
```

```
* DURING THE PROCESSING OF A SUBBAND SAMPLE STATE CAN CHANGE AS A
* RESULT OF AN INTERRUPT, SO DEFINE PATH=STATE/2
  LAC STATE,15
  SACH PATH
*
* CHECK IF DECODER IS ON (i.e. IN SYNCHRONIZATION)
  ZALS SWITCH
  BZ ACTSUB          SWITCH=0?
*
* DECODER IS NOT ON; OUTPUT A ZERO AND WAIT FOR NEXT CYCLE
  ZAC
  SACL OUTP
  B  WAITEV
*
* DECODER IS ON; ACTUAL SUBBAND DECODING
ACTSUB CALL PCMAQB
        CALL INVQMF
        B  WAITEV
*
-----
* INTERRUPT SERVICE ROUTINE
-----
*
INTRO  DINT
*
* SAVE STATUS OF TMS32010
  SST STATU
  MPY ONE          T REGISTER TO P REGISTER
  LDPK 1
  SAR 0,AR00
  SAR 1,AR01
  SACH ACH
  SACL ACL
  PAC
  SACL TREG
  LDPK 0
*
* STATE UPDATE
  LAC STATE
  ADD ONE
  AND MSTAT       STATE MOD 16
  SACL STATE
*
* RUN THROUGH DEMULTIPLEXER
  CALL DEMUX
*
* CREATE A 16 kHz OUTPUT SAMPLE OF THE TMS32010
* OUTPUT IS 8 kHz INVQMF TREE OUTPUT IF CYCLE IS ODD
* OUTPUT IS INTERPOLATION OF 8 kHz INVQMF TREE OUTPUT IF CYCLE IS EVEN
  LAC STATE
  AND ONE
  BZ INTPOL
  LAC OUTP        CYCLE IS ODD; OUTPUT IS INVQMF OUTPUT
  SACL BOUTP1
  OUT BOUTP2,2
  B  RSSTAT
```


APPENDIX G

Subroutine PCMAQB (decoder)


```

*****
#                               PCMAQB DECODING                               #
#                                                                           #
# THIS SUBROUTINE INCORPORATES:                                           #
#                               -BACKWARD STEPSIZE ADAPTION OF THE QUANTIZER #
#                               -2,3 or 4 BIT PCM DECODING OF A CODED SAMPLE #
#                               INTO A QUANTIZED SAMPLE OF THE SUBBAND BEING #
#                               PROCESSED                                  #
#                                                                           #
# INPUT:  FAW + 1 kHz SAMPLES (8 STREAMS)                                #
# OUTPUT: 1 kHz SAMPLES (8 STREAMS)                                       #
#-----#
# BIT ALLOCATION ACCORDING TO SSTRAT :                                     #
#                                                                           #
#                               1  2  3  4  5  6  7  8  -----#
#                               VOICED   4  4  3  2  2  0  0  0  #
#                               INTERMEDIATE 4  3  2  2  2  2  0  0  #
#                               UNVOICED  2  3  3  3  2  2  0  0  #
#-----#
# AUTHOR : ROLAND KLEUTERS                                               #
# DATE : 29-5-1987                                                       #
*****
#
PCMAQB  NOP                ENTRY POINT
#
#   READ IN FROM TABLE THE NUMBER OF CODE BITS FOR THE SUBBAND IN PROCESSION
#   LACK BAVD11
#   ADD  PATH
#   ADD  SSTRAT,3
#   TBLR NBIT                NBIT CONTAINS NUMBER OF CODE BITS
#
#   READ IN THE LOWEST AND HIGHEST POSSIBLE INPUT VALUE OF THE EXPONENT TABLE
#   LACK XI10
#   ADD  PATH,1
#   TBLR I0                  I0 CONTAINS LOWEST POSSIBLE INPUT OF EXP. TABLE
#   ADD  ONE
#   TBLR XI27                XI27 CONTAINS HIGHEST POSSIBLE INPUT OF EXP. TABLE
#
#   CHECK IF NUMBER OF CODE BITS (NBIT) GREATER THAN ZERO
#   LAC  NBIT
#   B6Z  READIN                DECODE?
#
#   NO DECODING HAS TO TAKE PLACE (NBIT=0)
#   ZAC                ZERO "CODED" RESULT; I(n)=0
#   SACL CALC1
#   LACK PCOD1
#   ADD  PATH
#   TBLW CALC1                I(n) SAVED IN PREVIOUS CODED SUBBAND SAMPLES TABLE
#   LACK DNI                  MINIMIZE STEPSIZE; d(n)=I0
#   ADD  PATH
#   TBLW X0                    d(n) SAVED IN TABLE
#   ZAC                ZERO "DECODED" RESULT; QUANTIZED SUBBAND SAMPLE:=0
#   SACL OUTPCM                QUANTIZED SUBBAND SAMPLE SAVED IN OUTPCM
#   RET
#

```

```

-----
* DETERMINE STEPSIZE i.e. DELTA(M) BY BACKWARD STEPSIZE ADAPTION
-----
*
* .....
* DETERMINE d(n)=LOG(DELTA(n))
* .....
*
* READ IN THE PREVIOUS CODED RESULT RESULT I(n-1) FROM TABLE
READIN LACK PC001
      ADD PATH
      TBLR CALC1          CALC1 CONTAINS THE PREVIOUS CODED RESULT I(n-1)
*
* DETERMINE m([I(n-1)])=LOG(M(I(n-1))) FROM TABLES
      LAC NBIT
      SUB ONE,1
      SACL CALC2
      LACK LOG22
      ADD CALC2,4
      ADD CALC1
      TBLR RES1          RES1 CONTAINS m([I(n-1)])=LOG(M([I(n-1)]))
*
* DETERMINE (1-GAMMA)*DC FROM TABLE
      LACK 06D1
      ADD PATH
      TBLR CALC1          CALC1 CONTAINS (1-GAMMA)*DC
*
* DETERMINE GAMMA*d(n-1) FROM TABLE
      LACK DN1
      ADD PATH
      TBLR RES0          RES0 CONTAINS d(n-1)
      LT GAMMA
      MPY RES0
      PAC
      SACH CALC2,1      CALC2 CONTAINS GAMMA*d(n-1)
*
* DETERMINE GAMMA*d(n-1)+m([I(n-1)])+(1-GAMMA)*DC
      ZALH CALC2
      ADD RES1,12
      ADD CALC1,8
      SACH CALC1          CALC1 CONTAINS GAMMA*d(n-1)+m([I(n-1)])+
                          (1-GAMMA)*DC, i.e d(n)
*
*
* CONTROL THE LIMITATION OF d(n) TO ITS RANGE AND SAVE d(n) IN TABLE
      LAC CALC1
      SUB X127
      B6Z SATDNH          d(n) GREATER THAN MAXIMUM POSSIBLE VALUE OF d(n)?
      LAC X0
      SUB CALC1
      B6Z SATDNL          d(n) LESS THAN MINIMUM POSSIBLE VALUE OF d(n)?
      B SAVEDN           NO SATURATION NEEDED
SATDNH LAC X127          SATURATE d(n) TO MAXIMUM POSSIBLE VALUE
      SACL CALC1
      B SAVEDN
SATDNL LAC X0           SATURATE d(n) TO MINIMUM POSSIBLE VALUE
      SACL CALC1

```

```
SAVEDN LACK DN1          SAVE d(n) IN TABLE
      ADD PATH
      TBLW CALC1
*
*.....
*                                     DETERMINE DELTA(n)=EXP(d(n))
*.....
*
* ROUND d(n) TO ONE OF THE ENTRY POINTS OF THE EXPONENT TABLE
      ZALH CALC1
      SUBH X0
      SACH CALC1          CALC1 CONTAINS d(n)-X0
      LT CALC1
      MPY STEP
      PAC
      SACH CALC1          CALC1 CONTAINS (d(n)-X0)*STEP
      LAC CALC1,11
      SACH CALC1          CALC1 CONTAINS INTEGER[(d(n)-X0)*STEP]=ENTRY POINT
*
* DETERMINE THE OUTPUT OF THE EXPONENT TABLE
      LACK CEXPO
      ADD CALC1
      TBLR CALC1          CALC1 CONTAINS OUTPUT OF EXPONENT TABLE=DELTA(n)
*
*-----
* DECODE PCMAQB DECODER INPUT (I(n)) INTO A QUANTIZED SUBBAND SAMPLE
*-----
*
* READ IN I(n) FROM TABLE OF CODED SUBBAND SAMPLES AND
* SAVE I(n) IN TABLE OF PREVIOUS CODED SUBBAND SAMPLES
      LACK COD1
      ADD PATH
      TBLR CALC2          CALC2 CONTAINS CODED RESULT I(n)
      LACK PCOD1
      ADD PATH
      TBLW CALC2          I(n) SAVED IN TABLE
*
* DETERMINE (I(n)+0.5)*DELTA(n)
      LAC CALC2,1
      ADD ONE
      SACL CALC2          CALC2 CONTAINS I(n)+0.5 (FORMAT=Q1)
      LT CALC2
      MPY CALC1
      PAC
*                                     ACCU CONTAINS (I(n)+0.5)*DELTA(n), i.e. THE
*                                     QUANTIZED SUBBAND SAMPLE (FORMAT=Q16,Q17 or Q20)
*
* CONTROL THE LIMITATION OF THE QUANTIZED SUBBAND SAMPLE TO ITS RANGE AND
* SAVE THE QUANTIZED SUBBAND SAMPLE (IN Q14,Q15 or Q18 FORMAT) IN OUTPCM
      SACL RES0          RES0=DDDD DDDD DDDD DD-- ; D=DATA BIT --DON'T CARE
      SACH RES1          RES1=SSSS SSSS SSSS SSSD ; S=SIGMBIT
      LAC RES1
      SUB ONE
      B6Z SATOPH          QUANTIZED SUBBAND SAMPLE GREATER THAN MAX. VALUE?
      ZAC
      SUB ONE,1
      SUB RES1
      B6Z SATOPL          QUANTIZED SUBBAND SAMPLE LESS THAN MIN. VALUE?
      LAC RES1,14        SAVE QUANTIZED SUBBAND SAMPLE WITHOUT SATURATION
```

```
SACL RES1          RES1=SD00 0000 0000 0000
LAC RES0,14
SACH RES0          RES0=--DD DDDD DDDD DDDD
LAC RES0
AND V3FFF          ACCU LOW=00DD DDDD DDDD DDDD
ADD RES1           ACCU LOW=SDDD DDDD DDDD DDDD (G:~,Q15,Q18 FORMAT)
SACL OUTPCM
RET
SATOPH LAC V3FFF   SATURATE QUANTIZED SUBBAND SAMPLE TO MAX. AND SAVE
ADD ONE,14
SACL OUTPCM
RET
SATOPL LAC ONE,15  SATURATE QUANTIZED SUBBAND SAMPLE TO MIN. AND SAVE
SACL OUTPCM
RET
```

*

*

PCMAQB DECODING DONE

=====

APPENDIX H
Subroutine INVQMF

```

*****
#                               INVQMF FILTERING                               #
#                                                                           #
# THIS SUBROUTINE INCORPORATES:                                           #
#                               -INVERSE QMF FILTER TREE                     #
#                               -8 kHz SPEECH SIGNAL RECONSTRUCTION          #
#                               -TAKE OVER OF NEW VOICING STRATEGY FOR      #
#                               PCM DECODING                                #
#                                                                           #
# INPUT:  1 kHz SAMPLES (8 STREAMS)                                       #
# OUTPUT: 8 kHz SAMPLES (1 STREAM)                                       #
#-----#
# SUBBAND PROCESSING : 1,8,4,5,2,7,3,6                                     #
#                                                                           #
# BIT ALLOCATION :                                                         #
#                               1  2  3  4  5  6  7  8  -----#
#                               VOICED   4  4  3  2  2  0  0  0  #
#                               INTERMEDIATE 4  3  2  2  2  2  0  0  #
#                               UNVOICED  2  3  3  3  2  2  0  0  #
#                                                                           #
# FAW's :                                                                 #
#   VOICED >76      INTERMEDIATE>96      UNVOICED>CC                    #
#-----#
# AUTHOR: ROLAND KLEUTERS                                               #
# DATE: 29-5-1987                                                       #
*****

```

```

#-----#
# MACRO'S                                                                 #
#-----#
#
#   MACRO TO REALIZE FILTERING IN LOW BRANCH OF FIRST INVQMF STAGE
FLTFLD $MACRO

```

- ZAC
- LT XF16
- MPY CF1
- LTD XF15
- MPY CF3
- LTD XF14
- MPY CF5
- LTD XF13
- MPY CF7
- LTD XF12
- MPY CF9
- LTD XF11
- MPY CF11
- LTD XF10
- MPY CF13
- LTD XF9
- MPY CF15
- LTD XFB
- MPY CF14
- LTD XF7
- MPY CF12
- LTD XF6
- MPY CF10
- LTD XF5
- MPY CF8
- LTD XF4

MPY CF6
LTD XF3
MPY CF4
LTD XF2
MPY CF2
LTA INPF
MPY CF0
APAC
SACL RES0
SACH RES1
ADD RES0
ADDH RES1
SACH OUTP
LAC INPF
SACL XF2
\$END

FILTER RESULT HAS TO BE AMPLIFIED WITH 2

8 kHz OUTPUT OF TMS32010

*

* MACRO TO REALIZE FILTERING IN HIGH BRANCH OF FIRST INVQMF STAGE

FLTFHI \$MACRO
ZAC
LT XF32
MPY CF0
LTD XF31
MPY CF2
LTD XF30
MPY CF4
LTD XF29
MPY CF6
LTD XF28
MPY CF8
LTD XF27
MPY CF10
LTD XF26
MPY CF12
LTD XF25
MPY CF14
LTD XF24
MPY CF15
LTD XF23
MPY CF13
LTD XF22
MPY CF11
LTD XF21
MPY CF9
LTD XF20
MPY CF7
LTD XF19
MPY CF5
LTD XF18
MPY CF3
LTA INPF
MPY CF1
APAC
SACL RES0
SACH RES1
ADD RES0
ADDH RES1
SACH OUTP

FILTER RESULT HAS TO BE AMPLIFIED WITH 2

8 kHz OUTPUT OF TMS32010

LAC INPF
SACL XF18
\$END

*
* MACRO TO REALIZE FILTERING IN LOW BRANCH OF SECOND INVQMF STAGE

FLTSLD \$MACRO
ZAC
LT XS8
MPY CS1
LTD XS7
MPY CS3
LTD XS6
MPY CS5
LTD XS5
MPY CS7
LTD XS4
MPY CS6
LTD XS3
MPY CS4
LTD XS2
MPY CS2
LTD INPS
MPY CS0
APAC
SACL RES0
SACH RES1
ADD RES0
ADDH RES1
SACH RES0
\$END

FILTER RESULT HAS TO BE AMPLIFIED WITH 2

4 kHz INPUT TO FIRST INVQMF STAGE

*
* MACRO TO REALIZE FILTERING IN HIGH BRANCH OF SECOND INVQMF STAGE

FLTSHI \$MACRO
ZAC
LT XS8
MPY CS0
LTD XS7
MPY CS2
LTD XS6
MPY CS4
LTD XS5
MPY CS6
LTD XS4
MPY CS7
LTD XS3
MPY CS5
LTD XS2
MPY CS3
LTD INPS
MPY CS1
APAC
SACL RES0
SACH RES1
ADD RES0
ADDH RES1
SACH RES0
\$END

FILTER RESULT HAS TO BE AMPLIFIED WITH 2

4 kHz INPUT TO FIRST INVQMF STAGE

*

* MACRO TO REALIZE FILTERING IN LOW BRANCH OF THIRD INVQMF STAGE

```
FLTTLO $MACRO
ZAC
LT XT6
MPY CT1
LTD XT5
MPY CT3
LTD XT4
MPY CT5
LTD XT3
MPY CT4
LTD XT2
MPY CT2
LTD INPT
MPY CT0
APAC FILTER RESULT HAS TO BE AMPLIFIED WITH 2
SACL RES0
SACH RES1
ADD RES0
ADDH RES1
SACH RES0 2 kHz INPUT TO SECOND INVQMF STAGE
$END
```

*

* MACRO TO REALIZE FILTERING IN HIGH BRANCH OF THIRD INVQMF STAGE

```
FLTTHI $MACRO
ZAC
LT XT6
MPY CT0
LTD XT5
MPY CT2
LTD XT4
MPY CT4
LTD XT3
MPY CT5
LTD XT2
MPY CT3
LTD INPT
MPY CT1
APAC FILTER RESULT HAS TO BE AMPLIFIED WITH 2
SACL RES0
SACH RES1
ADD RES0
ADDH RES1
SACH RES0 2 kHz INPUT TO SECOND INVQMF STAGE
$END
```

*

* MACRO TO LOAD DELAY VARIABLES OF SECOND INVQMF STAGE

```
LOADDS $MACRO X
LACK :X.S:
LARK 0,XS2
TBLR ++
ADD ONE
TBLR ++
ADD ONE
TBLR ++
ADD ONE
TBLR ++
ADD ONE
```

```
TBLR **
ADD ONE
TBLR **
ADD ONE
TBLR *
$END

*
*   MACRO TO SAVE DELAY VARIABLES OF SECOND INVQMF STAGE
SAVEDS $MACRO X
LACK :X.S:
LARK 0,XS2
TBLW **
ADD ONE
TBLW **
ADD ONE
TBLW **
ADD ONE
TBLW **
ADD ONE
TBLW **
ADD ONE
TBLW **
ADD ONE
TBLW **
ADD ONE
TBLW *
$END

*
*   MACRO TO LOAD DELAY VARIABLES OF THIRD INVQMF STAGE
LOADDT $MACRO X
LACK :X.S:
LARK 0,XT2
TBLR **
ADD ONE
TBLR **
ADD ONE
TBLR **
ADD ONE
TBLR **
ADD ONE
TBLR **
ADD ONE
TBLR *
$END

*
*   MACRO TO SAVE DELAY VARIABLES OF THIRD INVQMF STAGE
SAVEDT $MACRO X
LACK :X.S:
LARK 0,XT2
TBLW **
ADD ONE
TBLW **
ADD ONE
TBLW **
ADD ONE
TBLW **
ADD ONE
TBLW **
ADD ONE
TBLW *
$END

*
```

```
-----
* MACRO DEFINITION DONE; MAIN PART OF SUBROUTINE BEGINS
-----
*
INVQMF NOP          ENTRY POINT
*
*   SELECT PATH
    LACK 4
    SUB PATH
    BZ PATH6Z
    BZ HIGH4        SUBBAND 2
    ADD ONE
    BZ HIGH6        SUBBAND 7
    ADD ONE
    BZ HIGH5        SUBBAND 3
    B  HIGH7        SUBBAND 6
PATH6Z SUB ONE
    BZ LOW7         SUBBAND 5
    SUB ONE
    BZ LOW5         SUBBAND 4
    SUB ONE
    BZ LOW6         SUBBAND 8
*
*                               SUBBAND 1
*
-----
* THIRD INVQMF STAGE:
*
*           LOW4  0-500   Hz
*           HIGH4 500-1000 Hz
*           LOW5  1500-2000 Hz
*           HIGH5 1000-1500 Hz
*           LOW6  3500-4000 Hz; NOT CODED
*           HIGH6 3000-3500 Hz; NOT CODED
*           LOW7  2000-2500 Hz
*           HIGH7 2500-3000 Hz
*           TAKE OVER NEW VOICING STRATEGY
*
-----
*
LOW4  DMOV A4          DOUBLE BUFFER
*
*   INPUT CALCULATION
    ZALK B4
    SUBH C4
    SACH INPT          1 kHz INPUT TO FILTER
    LAC  OUTPCM
    SACL A4            BUFFERED 1 kHz INPUT TO THIRD INVQMF STAGE
*
*   LOAD DELAY VARIABLES OF LOW4
    LOADDT IXT2
*
*   FILTER 0-500 Hz
    FLTTLO
*
*   SAVE DELAY VARIABLES OF LOW4
    SAVEDT IXT2
*
*   BRANCH TO SECOND INVQMF STAGE
    B  LOW2
*
```

```
HIGH4  NOP                NO DOUBLE BUFFER
*
*  INPUT CALCULATION
*  ZALH C4
*  ADDH B4
*  SACH INPT              1 kHz INPUT TO FILTER
*  LAC OUTPCM
*  SACL C4              BUFFERED 1 kHz INPUT TO THIRD INVQMF STAGE
*
*  LOAD DELAY VARIABLES OF HIGH4
*  LOADDT XXT8
*
*  FILTER 500-1000 Hz
*  FLTTHI
*
*  SAVE DELAY VARIABLES OF HIGH4
*  SAVEDT XXT8
*
*  BRANCH TO SECOND INVQMF STAGE
*  B   LOW2
*
LOW5   DMOV A5            DOUBLE BUFFER
*
*  INPUT CALCULATION
*  LAC B5,15
*  SUB C5,15
*  SACH INPT              1 kHz INPUT TO FILTER
*  LAC OUTPCM
*  SACL A5              BUFFERED 1 kHz INPUT TO THIRD INVQMF STAGE
*
*  LOAD DELAY VARIABLES OF LOW5
*  LOADDT XXT14
*
*  FILTER 1500-2000 Hz
*  FLTTLO
*
*  SAVE DELAY VARIABLES OF LOW5
*  SAVEDT XXT14
*
*  BRANCH TO SECOND INVQMF STAGE
*  B   HIGH2
*
HIGH5  NOP                NO DOUBLE BUFFER
*
*  INPUT CALCULATION
*  LAC C5,15
*  ADD B5,15
*  SACH INPT              1 kHz INPUT TO FILTER
*  LAC OUTPCM
*  SACL C5              BUFFERED 1 kHz INPUT TO THIRD INVQMF STAGE
*
*  LOAD DELAY VARIABLES OF HIGH5
*  LOADDT XXT20
*
*  FILTER 1000-1500 Hz
*  FLTTHI
*
```

```
*   SAVE DELAY VARIABLES OF HIGH5
      SAVEDT XXT20
*
*   BRANCH TO SECOND INVQMF STAGE
      B   HIGH2
*
LOW6  NOP
*
*   THE 3500-4000 Hz BAND IS NOT CODED
*
*   ZERO 2 kHz INPUT TO SECOND INVQMF STAGE
      ZAC
      SACL RES0          2 kHz INPUT TO SECOND INVQMF STAGE
*
*   BRANCH TO SECOND INVQMF STAGE
      B   LOW3
*
HIGH6  NOP
*
*   THE 3000-3500 Hz BAND IS NOT CODED
*
*   ZERO 2 kHz INPUT TO SECOND INVQMF STAGE
      ZAC
      SACL RES0
*
*   BRANCH TO SECOND INVQMF STAGE
      B   LOW3
*
LOW7  DMOV A7          DOUBLE BUFFER
*
*   INPUT CALCULATION
      LAC B7,15
      SUB C7,15
      SACH INPT        1 kHz INPUT TO FILTER
      LAC OUTPCM
      SACL A7          BUFFERED 1 kHz INPUT TO THIRD INVQMF STAGE
*
*   LOAD DELAY VARIABLES OF LOW7
      LOADDT XXT26
*
*   FILTER 2000-2500 Hz
      FLTTLO
*
*   SAVE DELAY VARIABLES OF LOW7
      SAVEDT XXT26
*
*   BRANCH TO SECOND INVQMF STAGE
      B   HIGH3
*
HIGH7  NOP          NO DOUBLE BUFFER
*
*   INPUT CALCULATION
      LAC C7,15
      ADD B7,15
      SACH INPT        1 kHz INPUT TO FILTER
      LAC OUTPCM
      SACL C7          BUFFERED 1 kHz INPUT TO THIRD INVQMF STAGE
*
```

```
*   LOAD DELAY VARIABLES OF HIGH7
      LOADDT XXT32
*
*   FILTER 2500-3000 Hz
      FLTTHI
*
*   SAVE DELAY VARIABLES OF HIGH7
      SAVEDT XXT32
*
*   UPDATE FRAME AND TAKE OVER NEW VOICING STRATEGY IF END OF FRAME
      LAR 0,FRAME
      BANZ SFRAME
      LACK 7           END OF FRAME; PREPARE NEXT FRAME
      SACL FRAME
      DMOV STRAT      TAKE OVER NEW VOICING STRATEGY
      B   BRSEC
SFRAME SAR 0,FRAME   NOT END OF FRAME; PREPARE NEXT DATABLOCK
*
*   BRANCH TO SECOND INVQMF STAGE
BRSEC  B   HIGH3
*
*-----*
* SECOND INVQMF STAGE:
*           LOW2 0kHz-1kHz
*           HIGH2 1kHz-2kHz
*           LOW3 3kHz-4kHz
*           HIGH3 2kHz-3kHz
*-----*
*
LOW2   DMOV A2           DOUBLE BUFFER
*
*   INPUT CALCULATION
      ZALH B2
      SUBH C2
      SACH INPS          2 kHz INPUT TO FILTER
      LAC RES0
      SACL A2           BUFFERED 2 kHz INPUT TO SECOND INVQMF STAGE
*
*   LOAD DELAY VARIABLES OF LOW2
      LOADDS XXS2
*
*   FILTER 0-1 kHz
      FLTSLD
*
*   SAVE DELAY VARIABLES OF LOW2
      SAVEDS XXS2
*
*   BRANCH TO FIRST INVQMF STAGE
      B   LOW1
*
HIGH2  NOP              NO DOUBLE BUFFER
*
*   INPUT CALCULATION
      ZALH C2
      ADDH B2
      SACH INPS          2 kHz INPUT TO FILTER
      LAC RES0
      SACL C2           BUFFERED 2 kHz INPUT TO SECOND INVQMF STAGE
```

```
*
*   LOAD DELAY VARIABLES OF HIGH2
*   LOADS XXS10
*
*   FILTER 1-2 kHz
*   FLTSHI
*
*   SAVE DELAY VARIABLES OF HIGH2
*   SAVEDS XXS10
*
*   BRANCH TO FIRST INVQMF STAGE
*   B   LOW1
*
LOW3  DMOV A3           DOUBLE BUFFER
*
*   INPUT CALCULATION
*   LAC B3,15
*   SUB C3,15
*   SACH INPS          2 kHz INPUT TO FILTER
*   LAC RES0
*   SACL A3            BUFFERED 2 kHz INPUT TO SECOND INVQMF STAGE
*
*   LOAD DELAY VARIABLES OF LOW3
*   LOADS XXS18
*
*   FILTER 3-4 kHz
*   FLTSLO
*
*   SAVE DELAY VARIABLES OF LOW3
*   SAVEDS XXS18
*
*   BRANCH TO FIRST INVQMF STAGE
*   B   HIGH1
*
HIGH3  NOP             NO DOUBLE BUFFER
*
*   INPUT CALCULATION
*   LAC C3,15
*   ADD B3,15
*   SACH INPS          2 kHz INPUT TO FILTER
*   LAC RES0
*   SACL C3            BUFFERED 2 kHz INPUT TO SECOND INVQMF STAGE
*
*   LOAD DELAY VARIABLES OF HIGH3
*   LOADS XXS26
*
*   FILTER 2-3 kHz
*   FLTSHI
*
*   SAVE DELAY VARIABLES OF HIGH3
*   SAVEDS XXS26
*
*   BRANCH TO FIRST INVQMF STAGE
*   B   HIGH1
*
```

```
-----
* FIRST INVQMF STAGE:
*          LOW1  0kHz-2kHz
*          HIGH1 2kHz-4kHz
*-----
*
* LOW1  DMOV A1          DOUBLE BUFFER
*
* INPUT CALCULATION
*   ZALH B1
*   SUB  C1,14
*   SACH INPF,1        4 kHz INPUT TO FILTER
*   LAC  RES0
*   SACL A1            BUFFERED 4 kHz INPUT TO FIRST INVQMF STAGE
*
* FILTER 0-2 kHz
*   FLTFLD
*
* INVQMF DONE
*   RET
*
* HIGH1 NOP            NO DOUBLE BUFFER
*
* INPUT CALCULATION
*   LAC  C1,14
*   ADDH B1
*   SACH INPF,1        4 kHz INPUT TO FILTER
*   LAC  RES0
*   SACL C1            BUFFERED 4 kHz INPUT TO FIRST INVQMF STAGE
*
* FILTER 2-4 kHz
*   FLTFSH
*
* INVQMF DONE
*   RET
*
*
*          INVQMF FILTERING DONE
*-----
```


APPENDIX I

Subroutine DEMULTIPLEXING

```

*****
#                                     DEMULTIPLEXING                                     #
#                                                                                       #
# THIS SUBROUTINE INCORPORATES:                                                         #
#                                     -RECEIVING OF PCM CODED DATA AND                 #
#                                     -DEMULTIPLEXING                                 #
#                                                                                       #
# INPUT: 16 KBPS BITSTREAM                                                             #
# OUTPUT: FAW + 1 kHz SAMPLES (8 STREAMS)                                             #
#-----#
# SUBBAND PROCESSING : 1,8,4,5,2,7,3,6                                               #
#                                                                                       #
# BITALLOCATION :                                                                       #
#                                     -----1 2 3 4 5 6 7 8-----#
#                                     VOICED   4 4 3 2 2 0 0 0 #
#                                     INTERMEDIATE 4 3 2 2 2 2 0 0 #
#                                     UNVOICED  2 3 3 3 2 2 0 0 #
# FAW's :                                                                              #
#     VOICED >76      INTERMEDIATE >96      UNVOICED >CC #
#-----#
# AUTHOR: ROLAND KLEUTERS                                                             #
# DATA: 29-5-1987                                                                     #
#-----#
#
#-----#
# MACRO'S
#-----#
#
#     MACRO TO RECEIVE SIGNBIT AND EXTEND SIGN IN RECEIVE BUFFER
RECSB $MACRO
    IN  SIGIN,3      SIGNBIT RECEIVED VIA PIN 12 OF EXPANSION PORT
    LAC  ONE
    AND  SIGIN      SELECT SIGNBIT; SIGNBIT IS LSB OF ACCU
    SACL DATA      DATA=>0000 or >0001
    ZAC
    EXTEND SIGN IN RECEIVE BUFFER
    SUB  DATA
    SACL DATA      DATA=>0000 or >FFFF
    $END
#
#     MACRO TO RECEIVE ONE BIT
RECEIVE $MACRO
    LAC  DATA,1
    SACL DATA      PLACE RESERVED FOR NEW RECEIVE BIT
    IN  SIGIN,3      RECEIVE BIT RECEIVED VIA PIN 12 OF EXPANSION PORT
    LAC  ONE
    AND  SIGIN      SELECT RECEIVE BIT; RECEIVE BIT IS LSB OF ACCU
    XOR  DATA
    SACL DATA      RECEIVE BIT INSERTED INTO LSB OF DATA
    $END
#
#-----#
# MACRO DEFINITION DONE; MAIN PART OF SUBROUTINE BEGINS
#-----#
#
# DEMUX  NOP          ENTRY POINT
#

```

```
* CHECK IF DEMULTIPLEXER IS IN ALLIGNMENT (SWITCH NEQ >FF)
  LACK >FF
  XOR SWITCH
  BNZ RCMODE
*
*::::::::::::::::::::::::::::::::::::::::::::::::::::::::::::::::::::::::::
*                               DEMULTIPLEXER IS NOT IN ALLIGNMENT
*::::::::::::::::::::::::::::::::::::::::::::::::::::::::::::::::::::::::::
*
* FAW SEARCH OR 128 BIT (=FRAME) SKIP?
  LAC POINT
  BZ SKPBIT
*
* FAW SEARCH; INSERT ONE BIT IN SEARCH FOR A FAW
  RCEIVE
  LACK >FF
  AND DATA          SELECT 8 LSB's OF DATA
  SACL DATA         DATA CONTAINS 0000 0000 xxxx xxxx
*
* CHECK FAW
  LACK >76
  XOR DATA          VOICED?
  BZ VOIFND
  LACK >96
  XOR DATA          INTERMEDIATE?
  BZ INTFND
  LACK >CC
  XOR DATA          UNVOICED?
  BZ UNVFND
  RET                AT NEXT ENTRY AGAIN FAW SEARCH
*
* FAW FOUND, SET STRATEGY
VOIFND ZAC          STRATEGY=0 (VOICED)
      SACL STRAT
      B INSBIT
INTFND LACK 1       STRATEGY=1 (INTERMEDIATE)
      SACL STRAT
      B INSBIT
UNVFND LACK 2       STRATEGY=2 (UNVOICED)
      SACL STRAT
*
* INITIALIZATION FOR 128 BIT (=FRAME) SKIP
INSBIT LACK 127
      SACL COUNT          COUNT 128 BITS
      ZAC
      SACL POINT          NEXT ENTRY AT SKPBIT
      RET
*
* SKIP 128 BITS BEFORE SECOND FAW CHECK
SKPBIT RCEIVE
      LAR 0,COUNT
      BANZ CONTSK
*
```

* 128 BITS SKIPPED
ZAC
SACL ERROR NO ERROR ALLOWANCES
LACK >FF
AND DATA SELECT 8 LSB'S OF DATA
SACL DATA DATA CONTAINS 0000 0000 xxxx xxxx
B FAWCHK

* 128 BITS NOT YET SKIPPED
CONTSK SAR 0,COUNT
RET NEXT ENTRY AGAIN AT SKPBIT

* Frame Alignment Word CONFIRMATION

*
FAWCHK NOP
LACK >76
XDR DATA VOICED?
BZ CONF0
LACK >96
XDR DATA INTERMEDIATE?
BZ CONF1
LACK >CC
XDR DATA UNVOICED?
BZ CONF2
LAR 0,ERROR FAW NOT CONFIRMED; ERROR ALLOWED?
BANZ ERR ERROR ALLOWED?

*
* ERROR NOT ALLOWED; SO SWITCH TO MISALIGNMENT
LACK >FF
SACL SWITCH TURN OFF DECODER
SACL POINT
RET NEXT ENTRY AT FRAME SEARCH

*
* ERROR ALLOWED; DECREMENT NO. OF ERROR ALLOWANCES; STRATEGY UNCHANGED
ERR SAR 0,ERROR
B INISYN

*
* FAW CONFIRMED; SET STRATEGY AND NO. OF ERROR ALLOWANCES

CONF0 ZAC
SACL STRAT SET VOICED STRATEGY (=0)
LACK 4
SACL ERROR 4 FRAMES IN ERROR ALLOWANCE
B INISYN

CONF1 LACK 1
SACL STRAT SET INTERMEDIATE STRATEGY (=1)
LACK 4
SACL ERROR 4 FRAMES IN ERROR ALLOWANCE
B INISYN

CONF2 LACK 2
SACL STRAT SET UNVOICED STRATEGY (=2)
LACK 4
SACL ERROR 4 FRAMES IN ERROR ALLOWANCE

*

* INITIALIZE FOR A NEW FRAME

INISYN LACK >FF
SACL DEMO NEW ACCESS MARK
LACK 1
SACL BAND NEXT SUBBAND TO RECEIVE (0-500 Hz)

* SYNCHRONIZE DECODER IF NEEDED

LAC SWITCH
BZ INIOUT DECODER ALREADY ON, i.e. SYNCHRONIZATION NEEDED?
LACK 4
SACL SWITCH DELAY FOR DECODER TO SWITCH ON
LACK 13
SACL STATE SET STATE TO 13 (STATE=0 to 15)
LACK 7
SACL FRAME INITIALIZE DECODER STATUS
INIOUT RET NEXT ENTRY AT RECEIVE MODE

*
*:::
* DEMULTIPLEXER IS IN ALLIGNMENT; RECEIVE MODE
*:::
*

RCMODE NOP

* CHECK IF DECODER IS ON (SWITCH=0); IF NOT DECREMENT SWITCH
ZALS SWITCH
BZ FAWSUB
SUB ONE
SACL SWITCH

* SELECT FAW or ONE OF THE SUBBANDS FROM WHICH DATA HAS TO BE SENT

FAWSUB LACK 4
SUB BAND
BZ FWSB6Z
BZ NEXT4 SUBBAND 4 (1500-2000 Hz)
ADD ONE
BZ NEXT5 SUBBAND 5 (2000-2500 Hz)
B NEXT6 SUBBAND 6 (2500-3000 Hz)
FWSB6Z SUB ONE
BZ NEXT3 SUBBAND 3 (1000-1500 Hz)
SUB ONE
BZ NEXT2 SUBBAND 2 (500-1000 Hz)
SUB ONE
BZ NEXT1 SUBBAND 1 (0-500 Hz)
FAW

* RECEIVE FAW BITS

NFAW ZALS DEMO NEW ACCESS CHECK
BZ FSIG

* INITIALIZATION FOR A NEW FAW ACCESS

ZAC
SACL DEMO NO NEW ACCESS ANYMORE
SACL DATA CLEAR RECEIVE BUFFER
LACK 7
SACL COUNT 8 FAW BITS TO RECEIVE

```
*
* RECEIVE A FAW BIT AND INSERT IT INTO LSB OF DATA
FSIG RECEIVE
    LAR 0,COUNT
    BANZ CONT
*
* ALL FAW BITS RECEIVED; CHECK FAW
    B FAWCHK
*
* NOT ALL FAW BITS RECEIVED,PREPARE NEXT BIT
CONT SAR 0,COUNT
    RET
*
*-----*
* RECEIVE SUBBAND BITS
*-----*
*.....*
* A BIT OF SUBBAND 1 (0-500 Hz) HAS TO BE RECEIVED
*.....*
*
NEXT1 ZALS DEMO          NEW ACCESS CHECK
      BZ LOOP1
*
* INITIALIZATION FOR A NEW SUBBAND 1 ACCESS
      ZAC
      SACL DEMO          NO NEW ACCESS ANYMORE
      LACK 2
      XOR STRAT          UNVOICED?
      BNZ VOINI
*
* FURTHER INITIALIZATION FOR UNVOICED STRATEGY
      ZAC
      SACL COUNT          1 DATABIT TO RECEIVE
      B SIGN1
*
* FURTHER INITIALIZATION FOR VOICED AND INTERMEDIATE STRATEGY
VOINI LACK 2
      SACL COUNT          3 DATABITS TO RECEIVE
*
* RECEIVE THE SIGNBIT OF SUBBAND 1 AND EXTEND THE SIGN IN RECEIVE BUFFER
SIGN1 RECSB
      RET
*
* RECEIVE A SUBBAND 1 DATABIT AND INSERT IT INTO LSB OF DATA
LOOP1 RECEIVE
      LAR 0,COUNT
      BANZ CONT1
*
* ALL SUBBAND 1 BITS RECEIVED; UPDATE COD1 AND PREPARE NEXT SUBBAND
      LACK COD1
      TBLW DATA
      LACK >FF
      SACL DEMO          NEW ACCESS MARK
      LACK 4
      SACL BAND          NEXT SUBBAND TO RECEIVE (1500-2000 Hz)
      RET
*
```

```
* NOT ALL SUBBAND 1 BITS RECEIVED,PREPARE NEXT BIT
CONT1 SAR 0,COUNT
RET
*
*.....
* A BIT OF SUBBAND 2 (500-1000 Hz) HAS TO BE RECEIVED
*.....
*
NEXT2 ZALS DEMO NEW ACCESS CHECK
BZ LOOP2
*
* INITIALIZATION FOR A NEW SUBBAND 2 ACCESS
ZAC
SACL DEMO NO NEW ACCESS ANYMORE
LAC STRAT VOICED?
BNZ INUN2
*
* FURTHER INITIALIZATION FOR VOICED STRATEGY
LACK 2
SACL COUNT 3 DATABITS TO RECEIVE
B SIGN2
*
* FURTHER INITIALIZATION FOR INTERMEDIATE AND UNVOICED STRATEGY
INUN2 LACK 1
SACL COUNT 2 DATABITS TO RECEIVE
*
* RECEIVE THE SIGNBIT OF SUBBAND 2 AND EXTEND THE SIGN IN RECEIVE BUFFER
SIGN2 RECSB
RET
*
* RECEIVE A SUBBAND 2 DATABIT AND INSERT IT INTO LSB OF DATA
LOOP2 RCEIVE
LAR 0,COUNT
BANZ CONT2
*
* ALL SUBBAND 2 BITS RECEIVED; UPDATE COD2 AND PREPARE NEXT SUBBAND
LACK COD2
TBLW DATA
LACK >FF
SACL DEMO NEW ACCESS MARK
LACK 3
SACL BAND NEXT SUBBAND TO RECEIVE (1000-1500 Hz)
RET
*
* NOT ALL SUBBAND 2 BITS RECEIVED,PREPARE NEXT BIT
CONT2 SAR 0,COUNT
RET
*
*.....
* A BIT OF SUBBAND 3 (1000-1500 Hz) HAS TO BE RECEIVED
*.....
*
NEXT3 ZALS DEMO NEW ACCESS CHECK
BZ LOOP3
*
```

```
*   INITIALIZATION FOR A NEW SUBBAND 3 ACCESS
      ZAC
      SACL DEMO          NO. NEW ACCESS ANYMORE
      LACK 1
      XOR  STRAT        INTERMEDIATE?
      BNZ  VOUN3
*
*   FURTHER INITIALIZATION FOR INTERMEDIATE STRATEGY
      ZAC
      SACL COUNT        1 DATABIT TO RECEIVE
      B   SIGN3
*
*   FURTHER INITIALIZATION FOR VOICED AND UNVOICED STRATEGY
VOUN3  LACK 1
      SACL COUNT        2 DATABITS TO RECEIVE
*
*   RECEIVE THE SIGNBIT OF SUBBAND 3 AND EXTEND THE SIGN IN RECEIVE BUFFER
SIGN3  RECSB
      RET
*
*   RECEIVE A SUBBAND 3 DATABIT AND INSERT IT INTO LSB OF DATA
LOOP3  RECEIVE
      LAR  0,COUNT
      BANZ CONT3
*
*   ALL SUBBAND 3 BITS RECEIVED AND PERHAPS ALSO END OF 15 BIT DATABLOCK
*   UPDATE COD3 AND PREPARE NEXT THING TO RECEIVE
      LACK COD3
      TBLW DATA
      LACK >FF
      SACL DEMO          NEW ACCESS MARK
      LAC  STRAT        VOICED?
      BZ  NEXT6F        END OF 15 BIT DATABLOCK
      LACK 6
      SACL BAND          NEXT SUBBAND TO RECEIVE (2500-3000 Hz)
      RET
*
*   NOT ALL SUBBAND 3 BITS RECEIVED,PREPARE NEXT BIT
CONT3  SAR  0,COUNT
      RET
*
*   .....
*   A BIT OF SUBBAND 4 (1500-2000 Hz) HAS TO BE RECEIVED
*   .....
*
NEXT4  ZALS DEMO          NEW ACCESS CHECK
      BZ  LOOP4
*
*   INITIALIZATION FOR A NEW SUBBAND 4 ACCESS
      ZAC
      SACL DEMO          NO NEW ACCESS ANYMORE
      LACK 2
      XOR  STRAT        UNVOICED?
      BNZ  VOIN4
*
```



```
* FURTHER INITIALIZATION FOR UNVOICED STRATEGY
  LACK 1
  SACL COUNT      2 DATABITS TO RECEIVE
  B  SIGN4
*
* FURTHER INITIALIZATION FOR VOICED AND INTERMEDIATE STRATEGY
VOIN4 ZAC
  SACL COUNT      1 DATABIT TO RECEIVE
*
* RECEIVE THE SIGNBIT OF SUBBAND 4 AND EXTEND THE SIGN IN RECEIVE BUFFER
SIGN4 RECSB
  RET
*
* RECEIVE A SUBBAND 4 DATABIT AND INSERT IT INTO LSB OF DATA
LOOP4 RCEIVE
  LAR 0,COUNT
  BANZ CNT4
*
* ALL SUBBAND 4 BITS RECEIVED; UPDATE COD4 AND PREPARE NEXT SUBBAND
  LACK COD4
  TBLW DATA
  LACK >FF
  SACL DEMO      NEW ACCESS MARK
  LACK 5
  SACL BAND      NEXT SUBBAND TO RECEIVE (2000-2500 Hz)
  RET
*
* NOT ALL SUBBAND 4 BITS RECEIVED,PREPARE NEXT BIT
CONT4 SAR 0,COUNT
  RET
*
*.....
*          A BIT OF SUBBAND 5 (2000-2500 Hz) HAS TO BE RECEIVED
*.....
*
NEXT5 ZALS DEMO      NEW ACCESS CHECK
  BZ  LOOP5
*
* INITIALIZATION FOR A NEW SUBBAND 5 ACCESS
  ZAC
  SACL DEMO      NO NEW ACCESS ANYMORE
*
*          1 DATABIT TO RECEIVE FOR ALL STRATEGIES
*
* RECEIVE THE SIGNBIT OF SUBBAND 5 AND EXTEND THE SIGN IN RECEIVE BUFFER
SIGN5 RECSB
  RET
*
* RECEIVE THE SUBBAND 5 DATABIT AND INSERT IT INTO LSB OF DATA
LOOP5 RCEIVE
*
* ALL SUBBAND 5 BITS RECEIVED; UPDATE COD5 AND PREPARE NEXT SUBBAND
  LACK COD5
  TBLW DATA
  LACK >FF
  SACL DEMO      NEW ACCESS MARK
  LACK 2
  SACL BAND      NEXT SUBBAND TO RECEIVE (500-1000 Hz)
  RET
```

