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SIC-Based Detection With List and Lattice **Reduction for MIMO Channels**

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Abstract—To derive low-complexity multiple-input-multiple-output 5 (MIMO) detectors, we combine two complementary approaches, i.e., lat-6 tice reduction (LR) and list within the framework of the successive interfer-7 ence cancellation (SIC)-based detection. It is shown that the performance 8 of the proposed detector, which is called the SIC-based detector with list 9 and LR, can approach that of the maximum-likelihood (ML) detector with 10 a short list length. For example, the signal-to-noise ratio (SNR) loss of the 11 proposed detector, compared with that of the ML detector, is less than 1 dB 12 for a 4×4 MIMO system with 16-state quadrature amplitude modulation 13 (QAM) at a bit error rate (BER) of 10^{-3} with a list length of 8.

14 Index Terms-Lattice reduction (LR)-based detection, list detection, 15 successive interference cancellation (SIC), multiple-input-multiple-output 16 (MIMO) detection.

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I. INTRODUCTION

In wireless communications, it is well known that the channel 18 19 capacity can linearly increase with the number of antennas (provided 20 that the numbers of transmit and receive antennas are the same) [1], 21 [2]. Thus, to increase the channel capacity, the transmitter and receiver 22 can be equipped with multiple antennas, and the resulting channel 23 becomes a multiple-input-multiple-output (MIMO) channel. Various 24 space-time architectures for signal transmission over MIMO channels 25 are proposed to effectively exploit spatial and temporal diversity gain 26 in [3] and [4].

27 In general, since more symbols are transmitted in MIMO systems, 28 the detection complexity can be high. For example, the complexity of 29 maximum-likelihood (ML) detection exponentially increases with the 30 number of transmit antennas. Thus, various approaches are devised to 31 reduce the complexity. The successive interference cancellation (SIC) 32 approach is employed in [4]. The relation between SIC-based MIMO 33 detection and the decision feedback equalizer (DFE) is exploited in [5]. 34 In [6], the partial maximum a posteriori probability (MAP) principle is 35 derived to discuss the optimality of SIC-based detection. List detectors 36 are also considered for MIMO detection to obtain a soft decision in [7] 37 and [8] based on [9].

In [10], a lattice reduction (LR)-based MIMO detector used as 38 39 a low-complexity MIMO detector is first discussed. In [11], more 40 LR-based MIMO detectors are proposed. It is shown that the per-41 formance of LR-based MIMO detectors using minimum mean square 42 error (MMSE)-SIC approaches ML performance. An overview of LR-43 based detection can be found in [12]. In [13] and [14], it is shown that 44 LR-based detection can achieve full diversity. This is an important ob-45 servation as most low-complexity suboptimal MIMO detectors could

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Although the Lenstra-Lenstra-Lovasz (LLL) algorithm, which is 48 one of the LR algorithms, has a polynomial (average) complexity (for 49 a certain class of random channel matrices) [16], [17], the complexity 50 increases relatively rapidly with the number of basis vectors (or the 51 number of transmit antennas). Thus, for a large MIMO system, the 52 computational complexity of the LR-based detection would still be 53 high. To further reduce the complexity, we can decompose a large 54 MIMO detection problem into multiple small MIMO subdetection 55 problems with SIC, as in [6]. Due to SIC, this approach would suffer 56 from error propagation. To mitigate error propagation, the list detection 57 approach can be adopted. The resulting detector has low complexity as 58 the number of basis vectors in the subdetection problem is small. Due 59 to list detection, the proposed detector can enjoy the tradeoff between 60 complexity and performance, i.e., it has better mitigation against error 61 propagation as the list length increases at the expense of increasing 62 complexity. 63

II. SYSTEM MODEL 64

Suppose that there are K transmit antennas and N receive antennas. 65 The $N \times 1$ received signal vector **r** is given by 66

$$\mathbf{r} = \mathbf{H}\mathbf{s} + \mathbf{n} \tag{1}$$

where **H**, **s**, and **n** are the $N \times K$ channel matrix, $K \times 1$ transmitted 67 signal vector, and $N \times 1$ noise vector, respectively. We assume that 68 n is a zero-mean circular complex Gaussian random vector with 69 $E[\mathbf{nn}^{H}] = N_0 \mathbf{I}$. Let S denote the signal alphabet for symbols, i.e., 70 $s_k \in S$, where s_k is the kth element of s, and its size is denoted by M, 71 i.e., M = |S|. 72

We assume that $N \ge K$ and consider the QR factorization of the 73 channel matrix as H = QR, where Q is unitary, and R is upper 74 triangular. We have 75

$$\mathbf{x} = \mathbf{Q}^H \mathbf{r} = \mathbf{R}\mathbf{s} + \mathbf{Q}^H \mathbf{n}.$$
 (2)

Since the statistical properties of $\mathbf{Q}^{H}\mathbf{n}$ are identical to those of \mathbf{n} , 76 $\mathbf{Q}^{H}\mathbf{n}$ will be denoted by **n**. If N = K, there are no zero rows in **R**; 77 otherwise, the last N - K rows become zero. Thus, the last N - K 78 elements of x would be ignored for the detection if N > K. If there 79 is no risk of confusion, hereinafter, we assume that the sizes of x, R, 80 and **n** are $K \times 1$, $K \times K$, and $K \times 1$, respectively. 81

III. SIC-LIST-LR BASED DETECTION 82

The LR-based detectors in [10] and [11] have near-ML performance 83 with relatively low complexity. It is shown that those LR-based detec- 84 tors can achieve full diversity gain, just like the ML detector in [13] 85 and [14]. Unfortunately, however, the complexity of LR can rapidly 86 increase with the number of basis vectors, which implies that the 87 complexity of the LR-based detectors may not be reasonably low for a 88 large MIMO system. To avoid this problem, in this section, we propose 89 an SIC-list-LR-based detection method within the framework of the 90 partial MAP detection in [6]. The main idea of this method is to break 91 a high-dimensional MIMO detection problem into multiple lower 92 dimensional MIMO subdetection problems so that the complexity 93

94 associated with LR can be reduced. The notion of the partial MAP 95 detection [6] is applied to include multiple lower dimensional MIMO 96 subdetection problems, together with the list detection approach.

To perform the proposed LR and list-based detection, we consider 97 98 the partition of \mathbf{x} as follows:

$$\begin{bmatrix} \mathbf{x}_1 \\ \mathbf{x}_2 \end{bmatrix} = \begin{bmatrix} \mathbf{R}_1 & \mathbf{R}_3 \\ \mathbf{0} & \mathbf{R}_2 \end{bmatrix} \begin{bmatrix} \mathbf{s}_1 \\ \mathbf{s}_2 \end{bmatrix} + \begin{bmatrix} \mathbf{n}_1 \\ \mathbf{n}_2 \end{bmatrix}$$
(3)

99 where \mathbf{x}_i , \mathbf{s}_i , and \mathbf{n}_i are the $K_i \times 1$ *i*th subvectors of \mathbf{x} , \mathbf{s} , and \mathbf{n} , i =100 1, 2, respectively. Note that $K_1 + K_2 = K$. From (3), we can have two 101 lower dimensional MIMO subdetection problems to detect s_1 and s_2 . 102 It is straightforward to extend the partition into more than two groups. 103 However, for the sake of simplicity, we only consider the partition into 104 two groups, as in (3).

105 A. Algorithm Description

In the proposed SIC-list-LR-based detection, the subdetection of s_2 106 107 is carried out first using the LR-based detector. Then, a list of candidate 108 vectors of s_2 is generated. With the list of s_2 , the subdetection of s_1 is 109 performed with the LR-based detector. The candidate vector in the list 110 is used for the SIC to mitigate the interference from s_2 . The proposed 111 SIC-list-LR-based detection is summarized here.

112 S1) The LR-based detection of s_2 is performed with the received 113 signal \mathbf{x}_2 , i.e.,

$$\tilde{\mathbf{c}}_2 = \mathsf{LRDet}(\mathbf{x}_2) \tag{4}$$

114 where LRDet is the function of the LR detection operation,

which will be discussed in Section III-B, and \tilde{c}_2 is the estimated 115

116 vector of s_2 in the corresponding LR domain. Note that there is

- 117 no interference from s_1 in detecting s_2 .
- 118 S2) A list of candidate vectors in the LR domain is generated by

$$C_2 = \text{List}(\tilde{\mathbf{c}}_2)$$

(5)

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- where List is a function that chooses the Q closest vectors to 119
- 120 $\tilde{\mathbf{c}}_2(1 \leq Q \leq M^{K_2})$ in the LR domain. We will discuss the list generation in Section III-C. 121
- 122 S3) The list of candidates of s_2 , which is denoted by S_2 , can be converted from C_2 . For convenience, denote $S_2 =$ 123 $\{\tilde{\mathbf{s}}_{2}^{(1)}, \tilde{\mathbf{s}}_{2}^{(2)}, \dots, \tilde{\mathbf{s}}_{2}^{(Q)}\}.$ 124
- S4) Once S_2 is available, the LR-based detection of s_1 can be carried 125 out with SIC, i.e., 126

$$\mathbf{c}_{1}^{(q)} = \text{LRDet}\left(\mathbf{x}_{1} - \mathbf{R}_{3}\tilde{\mathbf{s}}_{2}^{(q)}\right)$$
(6)

where $\tilde{\mathbf{s}}_{2}^{(q)}$ is the *q*th decision vector of \mathbf{s}_{2} from list \mathcal{S}_{2} . 127

S5) Let $\tilde{\mathbf{s}}_1^{(q)}$ denote the signal vector corresponding to $\tilde{\mathbf{c}}_1^{(q)}$ in the LR 128 domain and $\tilde{\mathbf{s}}^{(q)} = \begin{bmatrix} (\tilde{\mathbf{s}}_1^{(q)})^T & (\tilde{\mathbf{s}}_2^{(q)})^T \end{bmatrix}^T$; the final decision of s 129 130 is found as

$$\tilde{\mathbf{s}} = \arg \min_{q=1,2,\dots,Q} \left\| \mathbf{x} - \mathbf{R} \tilde{\mathbf{s}}^{(q)} \right\|^2.$$
(7)

Note that a soft decision is also available from the list generated 131 132 in S5. There are Q candidate vectors for s, and they can be used to 133 approximate the log-likelihood ratio as a soft decision, as in [18]. In the 134 succeeding sections, we will explain the proposed detection in detail.

135 B. LR-Based Detection

In this section, we describe the LR-based detection used in steps S1) 136 137 and S4).

TABLE I SIGNALS AND PARAMETERS FOR THE LR-BASED DETECTION IN (4) AND (6)

Steps	У	Α	Z	ĉ	K_i
S1)	\mathbf{x}_2	\mathbf{R}_2	\mathbf{s}_2	$\tilde{\mathbf{c}}_2$	K_2
S4)	$\mathbf{x}_1 - \mathbf{R}_2 \hat{\mathbf{s}}_2^{(q)}$	\mathbf{R}_1	\mathbf{s}_1	$\tilde{\mathbf{c}}_1^{(q)}$	K_1



Fig. 1. List in different domains. (Left) C_2 in the LR domain, which is orthogonal and where the black dot represents \tilde{c}_2 . (Right) S_2 in the original domain, where the black dot represents \tilde{s}_2 .

Let ${\mathbb C}$ denote the set of complex integers or Gaussian integers 138 $\mathbb{C} = \mathbb{Z} + j\mathbb{Z}$, where \mathbb{Z} is the set of integers, and $j = \sqrt{-1}$. We assume 139 that $\{\alpha s + \beta | s \in S\} \subset \mathbb{C}$, where α and β are the scaling and shifting 140 coefficients, respectively. For example, for M-QAM, if $M = 2^{2m}$, 141 we have 142

$$S = \{s = a + jb | a, b \in \{\pm A, \pm 3A, \dots, \pm (2m - 1)A\}\}$$

v

where $A = \sqrt{(3E_s/2(M-1))}$, and $E_s = E[|s|^2]$ is the symbol 143 energy. Thus, $\alpha = 1/(2A)$, and $\beta = ((2m - 1)/2)(1 + j)$. Note that 144 the pair of α and β is not uniquely decided. 145 146

Consider the MIMO detection with the following signal:

$$\mathbf{r} = \mathbf{A}\mathbf{z} + \mathbf{v} \tag{8}$$

where **A** is a MIMO channel matrix, $\mathbf{z} \in S^{K_i}$ is the signal vector, and 147 **v** is a zero-mean Gaussian noise with $E[\mathbf{v}\mathbf{v}^H] = N_0\mathbf{I}$. We scale and 148 shift y as 149

$$\mathbf{d} = \alpha \mathbf{y} + \beta \mathbf{A} \mathbf{1}$$
$$= \mathbf{A}(\alpha \mathbf{z} + \beta \mathbf{1}) + \alpha \mathbf{v}$$
$$= \mathbf{A} \mathbf{b} + \alpha \mathbf{v}$$
(9)

where $\mathbf{1} = \begin{bmatrix} 1 & 1 & \cdots & 1 \end{bmatrix}^T$, and $\mathbf{b} = \alpha \mathbf{z} + \beta \mathbf{1} \in \mathbb{C}^{K_i}$. Let

$$\bar{\mathbf{A}} = \mathbf{A}\mathbf{U} \tag{10}$$

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where U is a unimodular matrix.¹ Using any LR algorithm, including 151 the LLL algorithm [16], we can find the value of U that makes the 152 column vectors of A shorter. It follows that 153

$$\mathbf{d} = \mathbf{A}\mathbf{U}\mathbf{U}^{-1}\mathbf{b} + \alpha\mathbf{v}$$
$$= \bar{\mathbf{A}}\mathbf{c} + \alpha\mathbf{v} \tag{11}$$

where $\mathbf{c} = \mathbf{U}^{-1}\mathbf{b}$. Note that, as the basis vectors are complex, we can 154 use complex LR algorithms [19] or convert a complex matrix into a 155 real matrix, as in [11]. The MMSE filter for estimating c is given by 156

$$\mathbf{W}_{\text{MMSE}} = \min_{\mathbf{W}} E\left[\left\| \mathbf{W}^{H} (\mathbf{d} - \bar{\mathbf{d}}) - (\mathbf{c} - \bar{\mathbf{c}}) \right\|^{2} \right]$$
$$= \left(\bar{\mathbf{A}} \text{cov}(\mathbf{c}) \bar{\mathbf{A}}^{H} + |\alpha|^{2} N_{0} \mathbf{I} \right)^{-1} \bar{\mathbf{A}} \text{cov}(\mathbf{c})$$
$$= \left(\mathbf{A} \mathbf{A}^{H} \alpha^{2} E_{s} + |\alpha|^{2} N_{0} \mathbf{I} \right)^{-1} \mathbf{A} \mathbf{U}^{-H} \alpha^{2} E_{s} \qquad (12)$$

¹A unimodular matrix is a square integer matrix with determinant ± 1 .



Fig. 2. Error probability with the list of c_2 for various list lengths.

157 where $\mathbf{\bar{d}} = E[\mathbf{d}] = \beta \mathbf{A} \mathbf{1}$, $\mathbf{\bar{c}} = E[\mathbf{c}] = \mathbf{U}^{-1}\beta \mathbf{1}$, and $\operatorname{cov}(\mathbf{c}) = 158 \ |\alpha|^2 \mathbf{U}^{-1} \mathbf{U}^{-H} E_s$ since

$$\operatorname{cov}(\mathbf{c}) = \operatorname{cov}(\mathbf{U}^{-1}\mathbf{b}) = \alpha^2 E_s \mathbf{I}.$$

159 The estimate of **c** is given by

$$\tilde{\mathbf{c}} = \bar{\mathbf{c}} + \mathbf{W}_{\text{MMSE}}^{H} (\mathbf{d} - \bar{\mathbf{d}}).$$
(13)

In Table I, the signals and parameters for the LR-based MMSE de-161 tection for each step are shown. Note that other approaches, including 162 the LR-based MMSE-SIC detector in [11] or non-LR-based detectors, 163 can also be used for subdetection.

164 C. List Generation in the LR Domain

165 To avoid or mitigate the error propagation, the use of a list of 166 candidate vectors of s_2 in detecting s_1 is crucial. Using the ML metric, 167 we can find the candidate vectors for the list S_2 . Let

$$f\left(\mathbf{r}\middle|\hat{\mathbf{s}}_{2}^{(1)}\right) \ge f\left(\mathbf{r}\middle|\hat{\mathbf{s}}_{2}^{(2)}\right) \ge \dots \ge f\left(\mathbf{r}\middle|\hat{\mathbf{s}}_{2}^{(M^{K_{2}})}\right)$$

168 where $f(\mathbf{r}|\mathbf{s})$ is the likelihood function of \mathbf{s} for a given \mathbf{r} , and $\hat{\mathbf{s}}_2^{(q)}$ is the 169 symbol vector that corresponds to the *q*th largest likelihood. With log-170 likelihood values, we can also find the candidate vectors as follows:

$$\left\|\mathbf{r} - \mathbf{R}_{2} \hat{\mathbf{s}}_{2}^{(1)}\right\|^{2} \leq \left\|\mathbf{r} - \mathbf{R}_{2} \hat{\mathbf{s}}_{2}^{(2)}\right\|^{2} \leq \cdots \leq \left\|\mathbf{r} - \mathbf{R}_{2} \hat{\mathbf{s}}_{2}^{(M^{K_{2}})}\right\|^{2}.$$

171 Therefore, the ML-based list becomes

$$S_2 = \left\{ \hat{\mathbf{s}}_2^{(1)}, \hat{\mathbf{s}}_2^{(2)}, \dots, \hat{\mathbf{s}}_2^{(Q)} \right\}.$$
 (14)

172 However, for each log-likelihood value, we need to perform a 173 matrix–vector multiplication. Thus, the resulting computational com-174 plexity could be high.

175 To avoid high computational complexity in generating the list, we 176 can find a suboptimal list in the LR domain that can be obtained with 177 a low complexity. Consider (9). According to Table I, let $\mathbf{A} = \mathbf{R}_2$, 178 $\mathbf{d} = \alpha \mathbf{x}_2 + \beta \mathbf{A} \mathbf{1}$, and $\mathbf{b} = \alpha \mathbf{s}_2 + \beta \mathbf{1}$. Then, from (10), we have

$$\|\mathbf{r} - \mathbf{R}_2 \mathbf{s}_2\| \propto \|\mathbf{d} - \mathbf{A}\mathbf{b}\| = \|\mathbf{d} - \mathbf{A}\mathbf{c}\|.$$
(15)



Fig. 3. BER performance of a 4×4 MIMO system with 16-QAM signaling.



Fig. 4. BER performance of a 4×4 MIMO system with 64-QAM signaling.

It is noteworthy that the metric on the right-hand side of (15) 179 is defined in the LR domain. Let \tilde{s}_2 be the signal vector in \mathcal{S}^{K_2} 180 corresponding to \tilde{c}_2 , and assume that \tilde{s}_2 is sufficiently close to $\hat{s}_2^{(1)}$. 181 Then, we can have $\mathbf{d} \simeq \mathbf{A} \tilde{\mathbf{c}}_2$. From this, the ML metric (ignoring a 182 scaling factor) for constructing the list in the LR domain becomes 183

$$\|\mathbf{d} - \bar{\mathbf{Ac}}\| \simeq \|\bar{\mathbf{A}}\tilde{\mathbf{c}}_2 - \bar{\mathbf{A}}\mathbf{c}\| = \|\tilde{\mathbf{c}}_2 - \mathbf{c}\|_{\bar{\mathbf{A}}^{H}\bar{\mathbf{A}}}$$
 (16)

where $\|\mathbf{x}\|_{\mathbf{A}} = \sqrt{\mathbf{x}^{H} \mathbf{A} \mathbf{x}}$ is a weighted norm. The list in the LR 184 domain becomes 185

$$\mathcal{C}_{2} = \left\{ \mathbf{c}_{2} \middle| \| \tilde{\mathbf{c}}_{2} - \mathbf{c} \|_{\bar{\mathbf{A}}^{H} \bar{\mathbf{A}}} < r_{\bar{\mathbf{A}}}(Q) \right\}$$
(17)

where $r_{\bar{\mathbf{A}}}(Q) > 0$ is the radius of an ellipsoid centered at $\tilde{\mathbf{c}}_2$, which 186 contains Q elements in the LR domain. If the column vectors of $\bar{\mathbf{A}}$ 187 or the basis vectors in the LR domain are orthogonal, $\bar{\mathbf{A}}^H \bar{\mathbf{A}}$ becomes 188 diagonal. Furthermore, if they have the same norm, $\bar{\mathbf{A}}^H \bar{\mathbf{A}} \propto \mathbf{I}$. Thus, 189 for nearly orthogonal basis vectors of almost equal norm, the list of \mathbf{c}_2 190 can be approximated as 191

$$\mathcal{C}_2 \simeq \tilde{\mathcal{C}}_2 = \left\{ \mathbf{c}_2 \middle| \| \tilde{\mathbf{c}}_2 - \mathbf{c} \| < r(Q) \right\}$$
(18)

TABLE II Impact of the Maximum Number of Column Swaps in the LR on BER (BER Rapidly Decreases With the Number of Column Swaps. Only a Few Number of Column Swaps Is Required)

E_b/N_0	$N_{\rm cs} = 1$	$N_{\rm cs} = 2$	$N_{\rm cs} = 3$	$N_{\rm cs} = 4$
6 dB	2.2472×10^{-2}	2.2459×10^{-2}	2.2459×10^{-2}	2.2459×10^{-2}
10 dB	1.9068×10^{-3}	1.9062×10^{-3}	1.9062×10^{-3}	1.9062×10^{-3}



Fig. 5. BER performance comparison of a 4 \times 4 MIMO system with 16-QAM signaling.

192 where r(Q) > 0 is the radius of a sphere centered at \tilde{c}_2 , which 193 contains Q elements. Since the LR provides a set of nearly orthogonal 194 basis vectors for the LR-based detection, we can see that the column 195 vectors in \bar{A} in (10) are nearly orthogonal, as shown in Fig. 1, with 196 a two-basis system. Let \tilde{S}_2 denote the list in the original domain 197 obtained from \tilde{C}_2 as in step S3). Since no matrix–vector multiplications 198 are required to generate \tilde{C}_2 or \tilde{S}_2 , we can use \tilde{S}_2 as the list in the 199 proposed detector to reduce computational complexity.

IV. SIMULATION RESULT

In this section, we present simulation results. We mainly focus 202 on the case of K = 4, particularly the case of $K_1 = K_2 = 2$. The 203 elements of **H** are independent zero-mean circular complex Gaussian 204 random variables with unit variance. This case is particularly interest-205 ing as the Gaussian reduction, which can find the two shortest vectors 206 in two-basis systems [10], [20], can be used for LR.

207 In the proposed SIC-list-LR-based detection, list length Q plays 208 a key role in the tradeoff between complexity and performance. In 209 general, it is desirable that the list has the true transmitted vector of c_2 . 210 If not, the proposed detector will have an incorrect decision. If Q in-211 creases, the error probability that $S_2(C_2)$ does not have the correct vec-212 tor of $s_2(c_2)$, which is denoted by $P_e(S_2)$ or $P_e(C_2)$, decreases. Error 213 probability $P_e(C_2)$ is considered for the MIMO system with 16-state 214 quadrature amplitude modulation (16-QAM), and N = K = 4. Sim-215 ulation results are shown in Fig. 2, where the error probabilities are 216 shown with two different lists in (17) and (18). As the list in (18) 217 is suboptimal, the performance is worse. However, this performance 218 degradation is not significant as the column vectors of $\overline{\mathbf{A}}$ are nearly 219 orthogonal.

220 The bit error rate (BER) performance of a 4×4 MIMO system 221 with 16-QAM signaling is shown in Fig. 3. In this case, a near-ML 222 performance can be achieved when $Q \ge 8$. For example, the signal-



Fig. 6. Complexity comparison of a 4 \times 4 MIMO system with 16-QAM signaling.

to-noise ratio (SNR) loss of the proposed detector, compared with that 223 of the ML detector, is less than 1 dB at a BER of 10^{-3} when Q = 8. 224

Fig. 4 shows the simulation results with 64-state quadrature ampli- 225 tude modulation (16-QAM). This result again confirms that the pro- 226 posed SIC-list-LR-based detector can provide a near-ML performance 227 with low complexity. At a BER of 10^{-3} , the SNR loss is less than 1 228 dB, compared with that of the ML detector when Q = 12. As the SNR 229 or E_b/N_0 increases, the SNR loss increases. However, by increasing 230 list length Q, this loss can be reduced as the list length can exploit 231 the tradeoff between performance and complexity. Note that a full 232 diversity may not be achieved by the proposed detector with a fixed list 233 length, as shown in Figs. 3 and 4. The relationship between diversity 234 order and list length needs to be investigated in the near future.

In the LR-based detection, since the number of column swaps in 236 the LR operation is not fixed, the complexity can vary from a channel 237 matrix to another. Thus, in practice, the maximum number of column 238 swaps can be fixed to limit the maximum complexity for two-basis 239 systems. It is shown in [10] that the two shortest vectors can be found 240 within two iterations for more than 99% of 2×2 random matrices 241 (of Rayleigh fading). However, when the number of column swaps 242 is limited, the basis vectors may not be properly reduced for some 243 channels, and the BER performance could be degraded because of it. 244 To see the impact of the maximum number of column swaps, a 245 simulation is considered with 16-QAM. Table II presents the BER 246 performance when the maximum number of column swaps $N_{\rm cs}$ is 247 limited. It is shown that the performance degradation is negligible, 248 even though $N_{\rm cs} = 1$.

For comparison purposes, we consider the BER performance of the 250 LR-based MMSE-SIC detector, which is the best LR-based detector 251 among the LR-based detectors proposed in [11]. The BER perfor- 252 mance results are shown in Fig. 5. It is shown that the proposed 253 detector can provide a performance that is better by about 1 dB than 254 the LR-based MMSE-SIC detector at a BER of 10^{-2} . Again, we 255

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256 can confirm that the combination of the LR-based detection and list 257 detection can improve the performance of the LR-based detector and 258 is an effective means to approach the ML performance.

For complexity comparison, we can take the upper bound on 259 260 the average number of LLL iterations in [17], which is given by 261 $\bar{N}_{cs} = K^2 \log K / (N - K + 1)$. (We ignore some minor terms to 262 simplify the comparison.) For 4 \times 4 MIMO channels, we have $\bar{N}_{cs} =$ 263 $K^2 \log K/(N-K+1) = 16 \log 4$ for the LR-based MMSE-SIC 264 detector and $\bar{N}_{cs} = 2(K/2)^2 \log((K/2)/((N/2) - (K/2) + 1)) =$ 265 8 log 2 for the proposed detector. This shows complexity reduction 266 by more than half in terms of LLL iterations. Note that the proposed 267 detector has additional complexity to build a list, which may offset the 268 complexity advantage of the proposed detector over conventional LR-269 based detectors [11]. To further see the complexity of each detector, 270 simulations are considered under the same environment, as shown in 271 Fig. 5. Fig. 6 shows the estimated flops using MATLAB execution time 272 that was obtained over all operations for each detector through simu-273 lations. The execution time is averaged over hundreds of thousands of 274 channel realizations. The Sphere Schnorr-Euchner algorithm [21] is 275 used for the ML decoding, whereas the LLL-reduced algorithm with 276 reduction factor $\delta = 3/4$ [16] is chosen for the LR-based MMSE-SIC 277 detector [11]. (This is the same as that in Fig. 5.) No limitation on the 278 number of iterations is imposed for any LR algorithm. The proposed 279 LR-based list detector clearly requires the lowest execution time. We 280 can also see that the execution time of the proposed detector is slightly 281 higher than half of the execution time of the LR-based MMSE-SIC 282 detector where the LLL-reduced algorithm is used.

V. CONCLUDING REMARK

In this paper, we have derived an SIC-list-LR-based detector for 85 MIMO detection using two complementary techniques, i.e., LR and 86 list detection, within a framework of SIC-based detection. It was 87 shown that the proposed detector has a near-ML performance with low 88 complexity. The list length plays a key role in the tradeoff between 89 performance and complexity. The performance is improved for a 290 longer list length, whereas the complexity increases with list length *Q*.

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