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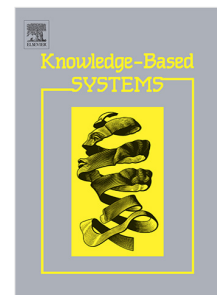
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Anjana Wijekoon: Conceptualization, Methodology, Software, Validation, Formal analysis, Investigation, Data Curation, Writing - Original Draft, Writing - Review & Editing, Visualization

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A User-centred Evaluation of DisCERN: Discovering Counterfactuals for Code Vulnerability Detection and Correction

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Abstract

Counterfactual explanations highlight *actionable knowledge* which helps to understand how a machine learning model outcome could be altered to a more favourable outcome. Understanding *actionable* corrections in source code analysis can be critical to proactively mitigate security attacks that are caused by known vulnerabilities. In this paper, we present the DisCERN explainer for discovering counterfactuals for code vulnerability correction. Given a vulnerable code segment, DisCERN finds counterfactual (i.e. non-vulnerable) code segments and recommends actionable corrections. DisCERN uses feature attribution knowledge to identify potentially vulnerable code statements. Subsequently, it applies a substitution-focused correction, suggesting suitable fixes by analysing the nearest-unlike neighbour. Overall, DisCERN aims to identify vulnerabilities and correct them while preserving both the code syntax and the original functionality of the code. A user study evaluated the utility of counterfactuals for vulnerability detection and correction compared to more commonly used feature attribution explainers. The study revealed that counterfactuals foster positive shifts in mental models, effectively guiding users toward making vulnerability corrections. Furthermore, counterfactuals significantly reduced the cognitive load when detecting and correcting vulnerabilities in complex code segments. Despite these benefits, the user study showed that feature attribution explanations are still more widely accepted than counterfactuals, possibly due to the greater familiarity with the former and the novelty of the latter. These findings encourage further research and development into counterfactual explanations, as they demonstrate the potential for acceptability over time among developers as a reliable resource for both coding and training.

Keywords: Counterfactual Explanations, Vulnerability Detection, Explainable AI

1. Introduction

Security attacks that exploit hidden software code flaws pose serious risks that compromise system performance and services. Therefore the ability to detect these vulnerabilities in a timely manner as well as being able to detect potential flaws is a desirable feature that can help to avoid financial and societal consequences. Application of AI for data-driven vulnerability detection has increased significantly in recent years [1, 2]. This is mainly due to the availability of large amounts of open-source code needed for training vulnerability detection models. Traditional classifiers such as SVM and Naive Bayes [3], as well as neural architectures for sequence modelling (e.g. LSTMs), have been successfully used for code vulnerability classification [4]. Given the structured textual nature of the data; these classifiers make use of text representation methods from information retrieval [3] as well as deep embedding techniques to represent software code [5].

Once vulnerabilities are detected or classified into flaw categories, the software needs to be fixed. Feature attribution methods enhance the transparency of AI model decisions by revealing the underlying reasoning for classifying a code segment as vulnerable. It assigns a weight to each token of the code which indicates how much it contributed to the AI model prediction (See examples in Figure 5). For example, authors of [1] used the feature activation map of their convolutional neural model to highlight parts of the code that contributed most to the AI model decision. Similarly authors of [6] used LIME to highlight the contribution of code tokens towards vulnerability. The methods introduced in this paper address a gap in the current approaches by focusing not only on identifying vulnerabilities but also on providing corrections as a solution. Here we demonstrate how research in counterfactual explanations can be conveniently adapted to generate code correction operators to guide the fixing of vulnerable code segments that are detected by a classification model.

Counterfactual Explanations for AI have accrued benefits from counterfactual thinking research from Psychology and GDPR guidelines for AI [7]. Counterfactuals reason with the inputs, the outputs, and the relationships between these to formulate a locally relevant explanation to convey how a

34 *better* or *more desirable* output (AI model decision) could have been achieved
35 by minimally changing the inputs. Questions concerning which part of the
36 input to modify and the appropriate methods for implementing such changes
37 to rectify code vulnerabilities are addressed in this paper. Here the input
38 is code segments and the proposed change relates to the code correction
39 operation. We present the DisCERN [8] algorithm, to locate the specific
40 area of vulnerability in a code segment, and to generate statement-level cor-
41 rections using substitution operations. In contrast to previous work where
42 DisCERN was employed for identifying substitutions using similarity calcu-
43 lations on tabular data, in this paper, substitutions are derived from code
44 snippets deemed similar but *non-vulnerable*. This is achieved by exploiting
45 similarity-driven pattern matching of pairs of code segments.

46 The utility of explanations in code vulnerability detection and correction
47 is best evaluated by the target users (i.e. developers). Accordingly, a user
48 study is performed to compare the effectiveness of counterfactuals from Dis-
49 CERN in comparison to feature attribution explanations from LIME. The
50 goal is to understand how counterfactuals and feature attributions differ in
51 the application of code vulnerability detection and correction in terms of
52 shaping mental models, affecting cognitive load and explanation goodness
53 and acceptability.

54 This paper makes the following contributions:

- 55 • introduces the DisCERN Counterfactual Explainer as a tool for code
56 vulnerability correction leveraging knowledge from feature attribution
57 explainers and pattern matching to make correction recommendations
58 (Section 4);
- 59 • demonstrates the generalisability of DisCERN across multiple program-
60 ming languages in terms of validity and sparsity metrics (Section 5);
61 and
- 62 • establishes the effectiveness of counterfactuals compared to feature at-
63 tribution explanations for vulnerability detection and correction in a
64 user study (Section 6).

65 The rest of the paper is organised as follows. Section 2 discusses the re-
66 lated work on vulnerability detection as a Machine Learning (ML) task and
67 correction from the view of XAI. The introduction of the NIST Datasets and
68 detection of code vulnerabilities using ML methods is presented in Section 3.

69 Section 4 presents the DisCERN algorithm which discovers counterfactuals
70 for vulnerable code segments and thereby guides the user to correct these vul-
71 nerabilities. The empirical evaluation and performance metrics with quanti-
72 tative and qualitative results are presented in Section 5. Section 6 presents
73 the user study that compares the utility of counterfactual vs feature attribu-
74 tion explanations. Finally, we draw conclusions in Section 7.

75 **2. Related Work**

76 *2.1. Code Vulnerability Detection*

77 The conventional approach to Code Vulnerability Detection (CVD) in-
78 volved software and security experts auditing a software system for potential
79 security defects, bugs and weaknesses all of which are referred to as vul-
80 nerabilities [9]. Automation of vulnerability detection of code is an active
81 applied research area where ML techniques are used for CVD [10, 11]. Early
82 ML methods for CVD focused on optimising feature extraction techniques
83 while neural network-based methods were used to learn semantic knowledge
84 from unstructured code to detect vulnerabilities [11]. Most recently, recur-
85 rent networks [12], graph neural networks [13] and transformer-based lan-
86 guage models [14, 15, 16] have been used for learning feature embeddings
87 from code for CVD. Many reviews in this research area provide comprehen-
88 sive overviews of ML techniques for CVD while emphasising the scarcity of
89 explainability approaches [10, 17]. XAI can be harnessed to support CVD in
90 multiple ways. For instance, it can help explain how the model works, iden-
91 tify the key features or variables that contribute to the detection process,
92 and provide insights into how to improve code and reduce vulnerabilities by
93 engaging humans in the loop. In this paper, we propose using the DisCERN
94 algorithm as a credible approach to address these issues.

95 *2.2. Code Vulnerability Correction*

96 The conventional approaches to providing users with corrective feedback
97 include rule-based [18, 19] and template-based approaches [20, 21]. Authors
98 of [18] proposed to pre-configure corrections for specific vulnerabilities and
99 reuse them as vulnerabilities are detected by their ensemble model in PHP
100 code. Similarly, authors of [19] use pre-configured vulnerability matching
101 rules and correction patterns for Java cryptography API code. Alternatively,
102 sequence-to-sequence models have been trained to generate corrections [22].

103 However, they are limited to a single programming language (C/C++) and
104 a vulnerability group (Buffer-overflow).

105 Our method is more closely related to work in [20] and [21] where the
106 methodology makes use of vulnerable and non-vulnerable code pairs to find
107 exemplar corrections. For each vulnerability in the code pair, they calcu-
108 late edit operations and cluster them to find correction patterns. Discovered
109 patterns are saved as templates to reuse on new vulnerable code segments.
110 Their method captures a wider variety of corrections by identifying multiple
111 correction patterns per vulnerability group. These methods share the same
112 challenge as DisCERN which is once a correction example (in DisCERN)
113 or a template(others) is found, how to adapt it to match the target code.
114 DisCERN addresses this by selecting the corrections from the nearest unlike
115 neighbour, which does not always guarantee perfect adaptation. Template-
116 based methods apply knowledge-intensive post-processing steps (such as cor-
117 recting variable names to match target code) that are not generalisable to
118 different languages and vulnerabilities.

119 The main difference between existing work and ours is that DisCERN is
120 generating the corrections to explain the prediction of an AI model (ex-
121 plaining the decision). Conversely, previous methods consider correction
122 generation to be an independent task and require a detection model that
123 classifies the exact vulnerability group. The difference is that DisCERN cor-
124 rections are guided by the knowledge encapsulated in the AI model such as
125 what features/tokens contributed to the decision. DisCERN is also not re-
126 liant on expert knowledge and heavily data-driven making it agnostic to the
127 detection-model and the programming-language. It also simplifies the task of
128 the detection AI model from a multi-class classification (up to 100+ classes)
129 problem to a binary-classification problem as the explainer does not require
130 the exact vulnerability group.


131 *2.3. Explainable AI in Vulnerability Detection*

132 Research literature and regulatory guidelines emphasise the necessity for
133 explanations of ML model decisions, as ML methods have increasingly be-
134 come more opaque and difficult to interpret [23, 24]. This applies to code
135 vulnerability detection and specifically towards prevention and or mitigation.
136 Feature attribution explainers have been explored as a way to pinpoint code
137 lines or segments that may have contributed to a *vulnerable* prediction by an
138 ML algorithm. Authors of [25] describe the design of a human-in-the-loop

139 XAI system for vulnerability mitigation, whereby model predictions are ex-
140 plained to forensic experts by way of feature attributions to enable them to
141 make necessary corrections. Authors of [26] explore the explanation needs
142 of target user groups of a code analyser to recognise two: a global expla-
143 nation where the common behaviours of the tool are explained; and a local
144 explanation where feature attribution explains why a specific code snippet is
145 predicted to be vulnerable. Both explanations are targeted towards a knowl-
146 edgeable audience of ML engineers. There are other works in similar areas
147 such as malware labelling in Android applications [27] and predicting phish-
148 ing URLs [28] that also make use of feature attribution explanations. Authors
149 of [6] used LIME to explain vulnerability detection in C/C++ code when us-
150 ing the Bidirectional LSTM model named VulDeePecker [12]. This paper
151 addresses a key gap in the literature by proposing the use of counterfactuals
152 not only for explaining detection but also for correcting vulnerabilities. Ac-
153 cordingly, [6] is the most directly linked previous work we compared against
154 DisCERN in our user study.

155 2.4. Explainable AI Techniques

156 Although there exists a broad range of explanation techniques and types [29]
157 our main emphasis is on factual and counterfactual explanations. The fac-
158 tual explanation often answers the “what” or “why” questions by providing
159 empirical evidence to support a particular AI model outcome based on the
160 input provided [30]. This evidence can take the form of feature attribution
161 where each input feature is assigned an attribution towards the outcome or
162 example-based explanations where nearest neighbours are used to support
163 the outcome. In contrast, counterfactuals answer “Why-not” or “How-to”
164 questions by formulating a hypothetical scenario that has a *more desirable*
165 outcome [30]. In code vulnerability detection and correction, a factual expla-
166 nation would highlight where the vulnerabilities exist within the code, while
167 a counterfactual explanation would help to demonstrate how to correct said
168 vulnerabilities. In this study, we investigate the use of the DisCERN algo-
169 rithm for discovering counterfactual explanations and evaluate its effective-
170 ness through a user study. The user study involves participants with varying
171 levels of expertise in code vulnerability detection and correction, allowing us
172 to assess the utility of the algorithm in a range of contexts.



```

173 public void method()
174 {
175     int data;
176     /* comment */
177     data = (new SecureRandom()).nextInt();
178     /* comment */
179     int array[] = { 0, 1, 2, 3, 4 };
180     /* comment */
181     if (data >= 0)
182     {
183         IO.writeLine(array[data]);
184     }
185     else
186     {
187         IO.writeLine("Array index out of bounds");
188     }
189 }

```

(a) Label: Vulnerable

```

173 public void method()
174 {
175     int data;
176     /* comment */
177     data = 2;
178     /* comment */
179     int array[] = { 0, 1, 2, 3, 4 };
180     /* comment */
181     if (data >= 0 && data < array.length)
182     {
183         IO.writeLine(array[data]);
184     }
185     else
186     {
187         IO.writeLine("Array index out of bounds");
188     }
189 }

```

(b) Label: Non-vulnerable

Figure 1: Pre-processed code segments from the Java dataset

173 3. Vulnerability Detection with NIST SAR Datasets

174 NIST Software Assurance Reference Dataset (SARD) Project promotes
 175 the detection and correction of known security flaws in programming code.
 176 The project maintains a publicly available repository of datasets from dif-
 177 ferent programming languages that are labelled for flaws and possible cor-
 178 rections. The flaws are standardised by the Common Weakness Enumera-
 179 tion (CWE) list which consists of software and hardware weaknesses. In this
 180 work, we consider three datasets in Java, C and C# programming languages
 181 from the NIST test suite ¹.

182 3.1. Preprocessing and Dataset Creation

183 In each dataset, code files are grouped under their CWE code and each file
 184 contains one or more functions (or methods in Java and C#). One function
 185 is *vulnerable* and often the remaining function is a proposed correction (i.e.
 186 non-vulnerable). We apply the following pre-processing steps to prepare each
 187 dataset for a binary-classification task:

- 188 1. Split functions in a file that are *vulnerable* and those *non-vulnerable*
 189 into individual data instances. An instance (i.e. function) was labelled

¹<https://samate.nist.gov/SARD/testsuite.php>

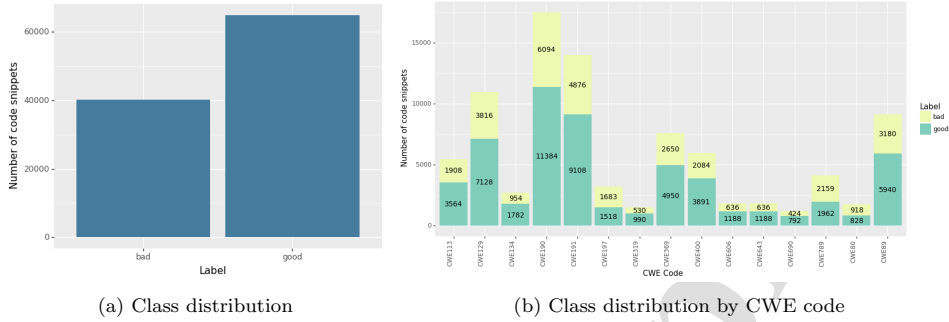


Figure 2: Java dataset statistics

190 *vulnerable* if it contains one or more comments that start with either
 191 *FLAW* or *POTENTIAL FLAW* and labelled as *non-vulnerable* if only
 192 contains comments that start with *FIX*.

- 193 2. Apply the following entity obfuscation steps to each function with the
 194 aim to prevent target leakage:
- 195 (a) replace all comments with `/*comment*/`; and
 - 196 (b) change all function signatures to `public void method()` (or lan-
 197 guage appropriate alternative).

198 Figure 1 presents two code segments from the Java dataset that were similar,
 199 one labelled as *vulnerable* and the other as *non-vulnerable*.

200 We present a detailed analysis of the class distribution of each dataset in
 201 Figures 2, 3 and 4. The left figure (Figure *a*) of each dataset shows that there
 202 are more *non-vulnerable* instances compared to *vulnerable* instances. Figure
 203 *b* on the right provides further analysis, examining the most frequent CWE
 204 codes (top 15) and the proportion of *vulnerable* and *non-vulnerable* instances
 205 for each code. Notably, there are no *non-vulnerable* examples for some CWE
 206 codes (example C# codes CWE313 and CWE94).

207 3.2. Vulnerability Classification

208 Code data can be seen as a text that follows grammar rules defined by the
 209 respective Compiler. The most common Machine Learning (ML) pipeline for
 210 classification with text data is to use a Tokenizer (t) to transform the text
 211 data into a vector representation and then apply a classification algorithm (f)
 212 to learn from labelled data. In this work, we consider several standard vec-
 213 tor representations and classifier combinations to compare the performance

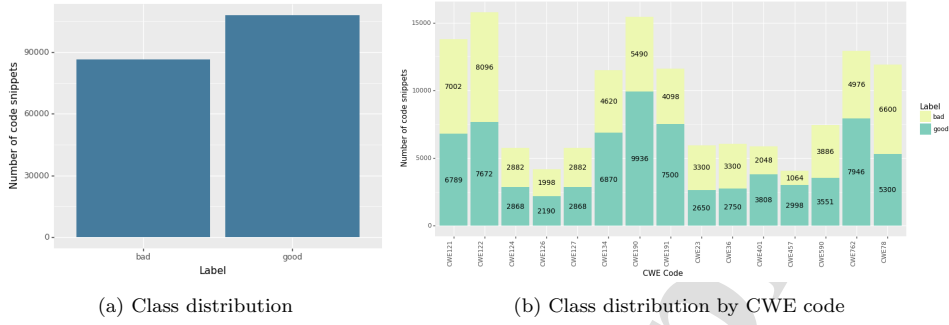


Figure 3: C dataset statistics

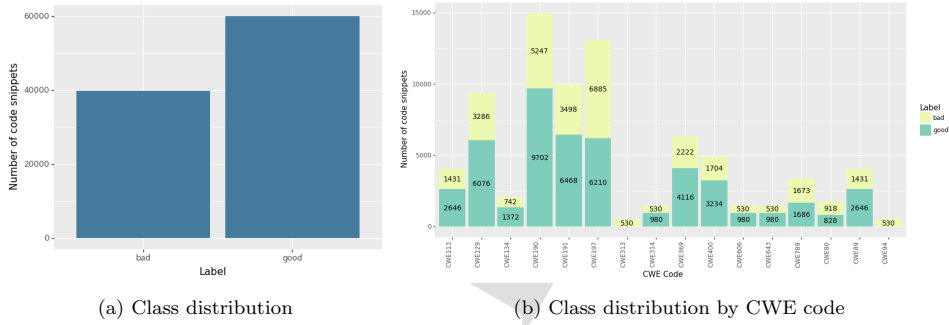


Figure 4: C# dataset statistics

214 of commonly used black box models that detect vulnerabilities in code seg-
 215 ments. We use 75/25 class stratified split to create 4 folds. For each fold,
 216 we train the model with 75% of the data and test with the remaining 25%.
 217 Table 1 presents the mean F1-score averaged across the four folds.

218 Overall we observe that *BoW + Random Forest* achieves the best perfor-
 219 mance for Java and C# datasets while CodeBERT classifier performs best
 220 for the C dataset. It is noteworthy that the contributions of this paper are
 221 model-agnostic, meaning that any combination of t and f should work with
 222 DisCERN, including the most recent encoders such as CodeBERT [31]. Ac-
 223 cordingly, the focus of the paper is not on identifying the best classification
 224 model, but rather to identify a model that performs well for experimental
 225 purposes. Accordingly, XAI evaluations in Section 5 used the *BoW + Ran-*
 226 *dom Forest* as the detection pipeline for all three datasets. This allowed for
 227 fairness and consistency across experiments and helped to observe the impact

Table 1: Classification Algorithms and Performance

Tokenizer (t)	Classifier (f)	Dataset		
		Java	C	C#
Tf-idf	Naive Bayes	0.7206	0.7284	0.7783
	kNN	0.9387	0.8457	0.9494
	SVM	0.9574	0.8839	0.9723
	Random Forest	0.9722	0.8734	0.9844
BoW	Random Forest	0.9761	0.8790	0.9889
CodeBERT-base Tokerniser	CodeBERT classifier	0.9469	0.9484	0.9880

228 of classification performance on the counterfactual generation.

229 4. DisCERN Counterfactuals for Vulnerability Detection and Cor- 230 rection

231 Code vulnerability detection decisions can be explained using different
232 types of explanations. As discussed in Section 2, it is commonly explained
233 using a factual explanation that uses feature attributions to explain the de-
234 cision and it is often targeted to knowledgeable users. Given a code segment
235 that is labelled *vulnerable*, a factual explanation will point to the part of the
236 code segment which led the AI model to label it as *vulnerable*. An exam-
237 ple factual explanation is shown in Figure 5a where text highlights indicate
238 *vulnerable* and *non-vulnerable* tokens in a Blue to Orange heat map scale.
239 For an expert, this type of explanation should be sufficient as they have the
240 knowledge to correct the vulnerability. In contrast, a counterfactual expla-
241 nation in Figure 5b will compare the given code segment with a similar yet
242 *non-vulnerable* code segment and make recommendations on how to correct
243 the vulnerability. Accordingly we argue that counterfactual explanations are
244 more informative for both expert and non-expert users, and in support of
245 this claim, we present the DisCERN algorithm for generating counterfactual
246 explanations specifically for code vulnerability correction.

247 4.1. Problem Definition

248 Consider a query code segment x , with m number of statements where the
249 i^{th} statement is denoted by s_i . If the vulnerability detection pipeline used to

```

public void method()
{
    int data = someMethod();

    /*comment*/
    data--;
    int result = (int)(data);

    IO.writeLine("result: " + result);
}

```

(a) Factual Explanation

```

1 1 public void method()
2 2 {
3 3     int data = someMethod();
3+ 4     int data;
4 4
5 5     /*comment*/
6 6     data--;
7 7     int result = (int)(data);
6+ 8     data = 2;
7+ 9
8+ 10     /*comment*/
9+ 11     int result = (int)(--data);
8 12
9 13     IO.writeLine("result: " + result);
10 14
11 15 }

```

(b) Counterfactual Explanation

Figure 5: Examples of feature attribution and counterfactual explanations

250 predict the code vulnerability consists of a Tokeniser, t , and a classification
 251 model, f , the decision predicted for x is y .

$$\begin{aligned}
 x &= [s_1, s_2, \dots, s_m] \\
 y &= f(t(x))
 \end{aligned}
 \tag{1}$$

252 For a given query x , having prediction, $y = \textit{vulnerable}$, there are four steps
 253 to discovering *non-vulnerable* counterfactuals with DisCERN:

- 254 1. find the Nearest Unlike Neighbour (NUN), \hat{x} from the train dataset \mathcal{X} ;
- 255 2. for each token z in x , find the attribution weights, using a feature attri-
 256 bution explainer (in this work we use LIME);
- 257 3. given a vulnerable token, z , in x , find statements pairs for correction,
 258 i.e. a list of statements in x and a list of candidate statements in \hat{x} as
 259 a potential vulnerability correction;
- 260 4. create an updated code segment, x' , by adapting the vulnerability cor-
 261 rection and check x' for decision change using the vulnerability detec-
 262 tion pipeline; and
- 263 5. repeat steps 3 and 4 until the detection pipeline predicts *non-vulnerable*.

264 Once the adapted code segment achieves the desired decision (i.e. non-
 265 vulnerable), it is identified as the counterfactual of the query. Next, we will
 266 explore each of these steps in detail.

267 *4.2. Finding the Nearest Unlike Neighbour*

268 Given a query x , the NUN, \hat{x} , is the nearest instance found in the train
 269 data with a different decision or label. In the context of counterfactual dis-
 270 covery, our query x is *vulnerable*. Selecting the NUN as the starting point, we
 271 expect 1) to minimise the actionable changes needed to flip the prediction i.e.
 272 with as few changes as possible; and 2) to preserve the original functionality
 273 of the code segment while correcting vulnerabilities. As in Equation 2, \hat{x} has
 274 n number of statements and the prediction is \hat{y} . Importantly, \hat{x} and x can
 275 have different number of statements (i.e. $n \neq m$) and should have different
 276 decisions (i.e. $\hat{y} \neq y$).

$$\begin{aligned} \hat{x} &= [\hat{s}_1, \hat{s}_2, \dots, \hat{s}_n] \\ \hat{y} &= f(t(\hat{x})) \mid \hat{y} \neq y \end{aligned} \quad (2)$$

277 To find the NUN by similarity, it is necessary to use an encoder (E) to
 278 transform code segments into a vector representation. This work used Code-
 279 BERT [31] to encode code segments. CodeBERT is based on the BERT [32]
 280 architecture and is state-of-the-art for natural language code search and code
 281 generation. It supports multiple programming languages making it most
 282 suited for this task. More specifically, we use the pre-trained weights from
 283 *codebert-base* shared in the Hugging Face repository² which is trained using
 284 bi-modal data (consisting of the code and its natural language description as
 285 two modalities) from CodeSearchNet.

286 Given a code query, x , the encoder E generates a vector representation,
 287 v , where the standard *codebert-base* encoding length, l , is 768 (Equation 3).
 288 From the train data set \mathcal{X} , we filter data instances for which $y_i \neq y$ and
 289 create the subset \mathcal{X}' . \mathcal{X}' represents all the *non-vulnerable* code segments
 290 that can be used to find a nearest-unlike-neighbour for x . Each data instance
 291 in \mathcal{X}' is encoded using the encoder E to obtain the set of vectors \mathcal{V}' . We
 292 use cosine similarity to find the NUN due to its robustness in comparing
 293 high-dimensional data, and its output range of -1 to 1 allows for a clear
 294 interpretation of similarity scores. We compute the cosine similarity between
 295 the query x , and any other instance, x_i as in Equation 3.

²<https://huggingface.co/microsoft/codebert-base>

$$v = E(x) \text{ and } v \in \mathbb{R}^l$$

$$\text{cosine}(x, x_i) = \frac{\sum_{j=1}^l v_{ij}v_j}{\sqrt{\sum_{j=1}^l v_{ij}^2} \sqrt{\sum_{j=1}^l v_j^2}} \quad (3)$$

296 Once the pair-wise similarity is computed (between x and each x_i in \mathcal{X}'),
 297 we select the train instance x_i from the pair with the highest similarity as
 298 the NUN of x . In the rest of this paper, this function is referred to as *nn*
 299 which given, query, x , train subset, \mathcal{X}' and the similarity metric, returns the
 300 NUN, \hat{x} .

301 4.3. Finding Feature Attribution Weights

302 Building upon counterfactual reasoning, DisCERN uses feature attribu-
 303 tion to reveal the most important code tokens or segments that contribute
 304 to an outcome of *vulnerable*. By selectively substituting only these segments,
 305 DisCERN can then identify the minimum changes needed to reverse that de-
 306 cision. The feature attribution explainers can provide the knowledge needed
 307 for identifying the code segments that need to be substituted. Accordingly,
 308 without loss of generalisability, this section describes the use of LIME ex-
 309 plainer to find feature attributions of the query to identify which parts of the
 310 code had contributed to it being labelled as *vulnerable*.

311 LIME is a model-agnostic feature attribution explainer that creates an
 312 interpretable model around a data instance to estimate how each feature
 313 contributed to the black-box model outcome [33]. LIME creates a set of
 314 perturbations within the query neighbourhood and labels them using the
 315 black-box model. This newly labelled dataset is used to create a linear in-
 316 terpretable model (e.g. a linear regression model). The resulting surrogate
 317 model is interpretable and only locally faithful to the black-box model (i.e.
 318 correctly classifies the input instance, but not all data instances outside its
 319 immediate neighbourhood). The new interpretable model is used to explain
 320 the black-box model outcome of the query. The explanation is formed by
 321 obtaining the linear model coefficients that indicate how each feature con-
 322 tributed to the outcome.

323 Our selection of LIME as the feature attribution explainer is motivated by
 324 the evidence from the literature. Authors of [6] proposed the use of LIME in
 325 the code vulnerability detection domain. Their evaluation demonstrated that
 326 the attributions correctly identify tokens that cause vulnerabilities. When

327 applying LIME in the context of code segment data, the *features* are the
 328 tokens identified by the Tokenizer, t , in the vulnerability detection pipeline.
 329 Accordingly, LIME can be used to understand the outcome of $f(t(x))$, by
 330 assigning an attribution, w , to each token which indicates how much the
 331 token contributes to the outcome.

$$LIME(x, t, f) \rightarrow \{w(z) \mid w(z) \in \mathbb{R}, z \in Z\} \quad (4)$$

332 If the vocabulary of code segments is Z , LIME assigns a weight w for each
 333 token $z \in Z$ (Equation 4). A positive weight ($w \geq 0$) indicates that the
 334 corresponding token contributes positively and a negative weight ($w < 0$)
 335 contributes negatively towards the outcome. We sort the weights using the
 336 partial order condition, \mathcal{R} , in Equation 5 to obtain the sorted list of tokens
 337 ordered from highest to lowest contribution towards the *vulnerable* outcome
 338 as Z' .

$$z_i \preceq_{\mathcal{R}} z_j \iff \mathcal{R} :: w(z_i) \geq w(z_j) \quad (5)$$

339 4.4. Substitution Algorithm

340 Given a token, z , in the query code segment, the goal of the substi-
 341 tution algorithm is to find a matching list of statements in the query and
 342 respective matches in the NUN to adapt the query such that it leads to a
 343 changed decision (i.e. vulnerable to non-vulnerable). To the best of our
 344 knowledge, existing feature attribution explainers identify the importance of
 345 tokens instead of code statements or segments. Instead of modifying the
 346 generic feature attribution explainers to operate at the statement level, we
 347 use a post-processing step to find the matching statements in the query that
 348 contains the token z , followed by a Pattern Matching (*pm*) algorithm to find
 349 matching lists of statements as presented in Algorithm 1. This allows for
 350 flexibility and compatibility of DisCERN with various existing attribution
 351 explainers.

352 We use a simple lookup function to identify all code statements (S')
 353 in the (adapted) query x' , that contain the token z (Line 1). The next
 354 steps (Lines 2- 5) of finding the vulnerable statements and their replace-
 355 ments from NUN are based on the hypothesis that if a statement s_j in S' is
 356 *vulnerable*, it must be *corrected* in the NUN. Accordingly, for a statement,
 357 s_j , in S' , first, we use a Pattern Matching algorithm to find a matching
 358 list of statements $s'_{[i:j]}$ from x' and $\hat{s}_{[v:w]}$ from \hat{x} . Here, the subscripts indi-
 359 cate the start and end indices of the list of statements and s_j is found within

Algorithm 1 substitute**Require:** $x' = [s'_1, s'_2, \dots, s'_m]$: (adapted) query**Require:** $\hat{x} = [\hat{s}_1, \hat{s}_2, \dots, \hat{s}_n]$: NUN as a list of statements**Require:** z : token in the query

```

1:  $S' \leftarrow [s \in x' \mid z \in s]$     ▷ find the list of statements in  $x'$  that include  $z$ 
2: for  $s_j \in S'$  do
3:    $s'_{[i:k]}, \hat{s}_{[v:w]} \leftarrow pm(s_j, [s'_1, s'_2, \dots, s'_m], [\hat{s}_1, \hat{s}_2, \dots, \hat{s}_n])$ 
4:    $c_j = \text{cosine}(E(s'_{[i:k]}), E(\hat{s}_{[v:w]}))$     ▷ calculate similarity
5: end for
6:  $(s', \hat{s}) \leftarrow \arg \max_{(s'_{[i:k]}, \hat{s}_{[v:w]})} c_j$     ▷ select maximum similarity pair
7:  $x' \leftarrow \text{replace}(x', s', \hat{s})$     ▷ replace  $s'$  in  $x'$  with  $\hat{s}$ 
8: return  $x'$ 

```

360 $s'_{[i:k]}$ (Line 3). A pattern-matching algorithm like the Gestalt Pattern Match-
 361 ing or Levenshtein Edit Distance can find the changes required to transform
 362 one string to another where the types of edits are *replace*, *delete* and *in-*
 363 *sert*. This paper used the Gestalt Pattern Matching algorithm implemented
 364 by `cdiffib` Python package³. We consider consecutive lists of statements
 365 rather than individual statements to preserve the grammatical structure of
 366 the programming language as closely as possible.

367 Next, we calculate the similarity between the two lists of statements using
 368 Cosine similarity (Line 4). Similar to Section 4.2 we use the *codebert-base*
 369 encoder to transform the list of statements to a vector representation and
 370 calculate the cosine similarity. Once we have all the $(s'_{[i:k]}, \hat{s}_{[v:w]})$ pairs, and
 371 their similarities, c_j , we select the pair, (s', \hat{s}) , that has the maximum sim-
 372 ilarity (Line 6). We assume a vulnerable code segment and its corrected
 373 counterpart are different yet carry some similarities. Accordingly, by select-
 374 ing the pair with the highest similarity from the remaining, we expect to
 375 discard those suggested by *pm* that are not vulnerability corrections. Note
 376 that *pm* only returns edit operations, not exact matches, hence the similarity
 377 score between a pair is always < 1 . Finally, in Line 7 we replace the list of
 378 statements s' in x' with the list of statement \hat{s} to return the new adapted
 379 query.

³<https://github.com/mduggan/cdiffib>

380 4.5. DisCERN Algorithm

Algorithm 2 DisCERN Algorithm

Require: $x = [s_1, s_2, \dots, s_m]$: query as a list of statements**Require:** $f(t(\cdot))$: vulnerability detection pipeline**Require:** sim : similarity metric, default is cosine similarity**Require:** \mathcal{X} : train dataset**Require:** $y = f(t(x))$: black-box prediction for the query

```

1:  $\mathcal{X}' \leftarrow \{x_i \in \mathcal{X} \mid y_i \neq y\}$  ▷ filter the train dataset
2:  $\hat{x} \leftarrow nm(x, \mathcal{X}', sim)$  ▷ find the NUN
3:  $\{w(z)\} \leftarrow LIME(x, t, f)$  ▷ feature attributions
4:  $Z' \leftarrow \mathcal{R}(\{w(z)\})$  ▷ tokens sorted by  $\mathcal{R}$ 
5: Initialise  $x' = x$  and  $y' = y$ 
6: for  $z \in Z'$  do ▷ for each token in the sorted list
7:    $x' \leftarrow substitute(x', \hat{x}, z)$  ▷ Algorithm 1
8:    $y' = f(t(x'))$  ▷ predict decision for the adapted query  $x'$ 
9:   if  $y' \neq y$  then ▷ check if the decision is changed
10:     Break ▷ stop substitutions if decision is changed
11:   end if
12: end for
13: return  $x'$  ▷ return the adapted query as the counterfactual

```

381 DisCERN (Algorithm 2) brings together Sections 4.2 to 4.4 to discover
382 counterfactuals for vulnerable code. Given the query x , and the train dataset
383 \mathcal{X} , in Lines 1 and 2 we find the NUN as discussed in Section 4.2. Next,
384 we find the LIME feature weights for the query and sort it to obtain the
385 list of tokens that indicate which parts of the code contributed to the cur-
386 rent decision (Line 3 and 4, Section 4.3). We iterate over the sorted list of
387 tokens where for each token we find corresponding statements and substi-
388 tutions (from Algorithm 1) until the prediction is changed (Line 8). Here
389 the prediction for the adapted query x' is obtained using the original clas-
390 sification pipeline $f(t(\cdot))$. The iteration is terminated when a prediction is
391 changed and the algorithm returns the adapted query x' as the counterfac-
392 tual for the query x . Compared to DisCERN for tabular data [8] the key
393 novelty is the substitution algorithm that aims to preserve programme lan-
394 guage syntax and original functionality while correcting the vulnerabilities.
395 However, the outcome of, the substitution algorithm is dependent on the

396 Nearest-Unlike-Neighbour and does not always guarantee to find a counter-
 397 factual from the NUN. Accordingly, in the worst-case scenario, DisCERN
 398 iterates through all tokens in Z' and may fail to lead to a desirable deci-
 399 sion change (of *non-vulnerable*) even after all corrections are actioned on the
 400 query.

401 5. Evaluation

402 This section presents the evaluation of the counterfactual DisCERN al-
 403 gorithm for vulnerable code correction. To the best of our knowledge, there
 404 are no existing algorithms in the literature for counterfactual discovery in
 405 the code vulnerability correction domain to compare performance with other
 406 methods.

407 5.1. Performance Metrics

408 DisCERN algorithm is evaluated using the three NIST datasets (Sec-
 409 tion 3); in each dataset, we only use *vulnerable* test data instances for the
 410 XAI evaluations. The following metrics are used to measure the performance.

- 411 • **Validity** measures the percentage of data for which the algorithm suc-
 412 cessfully finds a counterfactual [34, 35, 8]. At this stage, the require-
 413 ment for a counterfactual discovered by an algorithm is to achieve a
 414 *positive* change of decision⁴. Given the set of test instances that were
 415 predicted *vulnerable* are X_v , and the subset for which the algorithm
 416 found a counterfactual is X_v^c , the validity is calculated as in Equation 6.
 417 A higher percentage of validity is desirable.

$$Validity = \frac{|X_v^c|}{|X_v|} \times 100 \quad (6)$$

- 418 • **Sparsity** measures the mean number of statements that were changed (i.e.
 419 cost) for a change in decision [34, 35, 8]. Given the cost for each test in-
 420 stance in X_v^c is $[r_1, r_2, \dots, r_N]$, where $N = |X_v^c|$, the sparsity is calculated
 421 as in Equation 7. In Algorithm 1, the number of statements changed for
 422 *replace*, *delete* and *insert* operations are calculated as $\max(k - i, w - v)$,

⁴A more stringent metric would be to evaluate if the change conforms to grammar rules of the Language Compiler, which we will explore in future work.

423 $k - i$ and $w - v$ respectively. As such, the cost of a test instance is
 424 determined by aggregating the number of statement changes that cor-
 425 respond to the applied operations. In other domains, lower sparsity
 426 is preferred, however, in this domain, we hypothesise sparsity is not
 427 directly correlated to the algorithm performance as a vulnerability cor-
 428 rection could require adding more statements. This will be discussed
 429 further with empirical results in Section 5.2.

$$Sparsity = \frac{1}{N} \sum_{j=1}^N r_j \quad (7)$$

430 There are other metrics used in counterfactual evaluations such as prox-
 431 imity (measures the difference between the original and the substitution code
 432 segments) [34, 35, 8] and diversity (measures the difference between multiple
 433 counterfactuals) [34] which we did not find to be transferable to the code
 434 vulnerability correction domain.

435 5.2. Results

436 Table 2 presents the performance evaluation results of DisCERN using the
 437 three NIST datasets. In addition to performance metrics, we also measure
 438 the mean number of statements in a query, nearest-unlike-neighbour and
 439 counterfactual which we found useful when discussing the performance of
 440 DisCERN.

Table 2: Validity and Sparsity of DisCERN

Dataset	Validity (%)	Sparsity	Mean no of statements in the		
			Query	NUN	CF
Java	96.49	13.88	44.62	51.81	50.93
C	85.50	8.40	24.78	28.26	26.08
C#	97.55	13.16	27.67	33.96	33.44

441 We observe that the validity is consistently below 100% across all datasets.
 442 The validity for the C dataset is significantly lower which means the C dataset
 443 queries were not able to find counterfactuals using DisCERN. This can be
 444 linked to a high ($\sim 21\%$) classification error seen in the vulnerability de-
 445 tection pipeline. For example, the query can be misclassified as vulnerable

446 or the adapted query can be continuously misclassified as *vulnerable*. It is
447 further validated by the Java and C# datasets showing validity consistent
448 with their classification pipeline performance.

449 Sparsity is measured as the number of changes that were required to
450 get the decision changed from *vulnerable* to *non-vulnerable*. Considering the
451 mean number of statements in the query (column 4), Java and C datasets
452 make less number of changes compared to C#. It is noteworthy that these
453 changes include *deletion* operations, thus it is not an indication of the length
454 of the counterfactual. When generating counterfactuals for tabular data, a
455 common goal is to minimise sparsity. However, when discovering counter-
456 factuals for correcting code vulnerabilities we argue that lower sparsity is
457 not always desirable. In general, correcting vulnerabilities can be costly; for
458 example in Java, adding a *try-catch-finally* block surrounding a vulnerable
459 statement can add up to 4-10 lines based on the formatting styles (Allman
460 vs K&R).

461 Further analysis of the number of statements between NUN and the coun-
462 terfactual shows the effectiveness of the DisCERN algorithm. The mean
463 number of statements in a CF is consistently lower than that in the NUN
464 indicating that DisCERN is in fact finding meaningful corrections instead
465 of completely converting the query into its NUN. The consistently higher
466 number of statements in CF compared to the Query further indicates the
467 increased cost of correcting code vulnerabilities.

468 5.3. Qualitative Analysis

469 While DisCERN aims to maintain syntactic integrity and preserve the
470 originally intended code functionality, sparsity, and validity metrics do not
471 specifically measure these aspects. As a result, we examined a selection of
472 the generated counterfactuals to determine whether the proposed code adap-
473 tations can effectively address code vulnerabilities and to what extent they
474 implement reasonable modifications without compromising functionality.

475 Consider the two illustrative Java code examples in Figures 6a and 6b
476 which were counterfactuals discovered by the DisCERN algorithm. In each
477 figure, the first two columns indicate the line numbers of the query and the
478 counterfactual; the third column uses addition and subtraction signs to in-
479 dicate adaptation operations. In example 1, a replacement is proposed (i.e.
480 replace query lines 5-6 with NUN lines 5-12). With Example 2, the coun-
481 terfactual proposes an insertion (i.e. insert new lines 4-6) and a replacement
482 (i.e. replace query lines 6-7 with NUN lines 9-13). Both sets of adaptations

<pre> 1 1 public void method() 2 2 { 3 3 float data = dataArray[2]; 4 4 /*comment*/ 5 5 int result = (int)(100.0 % data) 6 6 IO.writeLine(""); 7 7 if (Math.abs(data) > 0.000001) 8 8 { 9 9 int result = (int)(100.0 % data) 10 10 IO.writeLine(""); 11 11 } 12 12 else{ 13 13 IO.writeLine(""); 14 14 } 15 15 }</pre>	<pre> 1 1 public void method() 2 2 { 3 3 int data = dataArray[2]; 4 4 /*comment*/ 5 5 int array[] = null; 6 6 /*comment*/ 7 7 if (data > 0) 8 8 { 9 9 /*comment*/ 10 10 IO.writeLine((short) data); 11 11 array = new int[data]; 12 12 } 13 13 else 14 14 { 15 15 IO.writeLine(""); 16 16 } 17 17 }</pre>
(a) Example 1: Successful Adaptation	(b) Example 2: Unsuccessful Adaptation

Figure 6: DisCERN Counterfactual Examples

483 have maintained the grammatical structure of the Java language, however,
484 Example 1 is better at preserving functionality, because it ensures that the
485 original functionality of writing an empty line (originally line 6) even after
486 having introduced an *if* condition. In Example 2, DisCERN fails to preserve
487 the intended functionality in the original query line 7 (by failing to treat *data*
488 as an array).

489 Both examples corroborate findings in Table 2 that code vulnerability
490 correction can increase sparsity due to the insertion of additional statements.
491 Overall, both evaluations indicate that DisCERN is a promising approach
492 to discovering counterfactuals, however, to ensure comprehensive validity,
493 further adaptation heuristics are needed to verify counterfactuals maintain
494 the original functionality (e.g., apply unit testing if available).

495 6. User Evaluation

496 The primary objective of this user study is to assess the effectiveness
497 of factual and counterfactual explainers in addressing code vulnerabilities,
498 specifically examining their utility for both experienced and novice develop-
499 ers. While existing literature [25, 26, 27] highlights a focus on factual ex-
500 planations (such as feature attributions) for knowledgeable users in the XAI
501 research, our hypothesis posits that counterfactual explanations may prove
502 more informative for both skilled and trainee developers aiming to correct
503 code vulnerabilities. Table 3 presents the user study protocol; enumeration

504 indicates the order in which the questions were presented; Green colour indi-
 505 cates content presented to the participant (code segment or explanation) and
 506 the protocol is grouped by different intents (Blue). The questionnaire was
 507 prepared to capture users' mental models before and after receiving explana-
 508 tions, as well as to evaluate the quality and acceptability of the explanations
 509 provided by the system for detecting and correcting code vulnerabilities.

Table 3: User Study Protocol

Present code snippet	
A priori mental model for detecting code vulnerabilities	
Q1. Do you think the code snippet contains code vulnerabilities?	<i>Yes, No, Maybe</i>
A priori mental model for correcting code vulnerabilities	
Q2. If you answered yes, which lines would you change to correct code vulnerabilities?	<i>Free text</i>
Q3. If you listed any lines, why do you think these lines contain code vulnerabilities?	<i>Free text</i>
Present explanation (annotated or modified code snippet)	
A posterior mental model for correcting code vulnerabilities	
Q4. After seeing the explanation, which lines would you change to correct code vulnerabilities?	<i>Free text</i>
Q5. If you changed your answer from before viewing the explanation, please mention why?	<i>Free text</i>
Measure goodness of the explanation for detection and correction	
Q6. Did the explanation help you detect vulnerabilities?	<i>Yes, No</i>
Q7. Did the explanation help you to identify the lines you would change to correct code vulnerabilities?	<i>Yes, No</i>
Measure acceptability of the explanation	
Q8. Did the explainer correctly annotate the parts of the code that contain vulnerabilities?	<i>Yes, No, Partially</i>

510 The questionnaire was repeated with three different code snippets of dif-
 511 ferent lengths (11, 33 and 21 lines of code) to minimise bias. Snippets were
 512 selected from the Java dataset over *C* and *C#* languages considering the
 513 wider usage and familiarity within the target user group. All snippets con-
 514 tained a variant of the CWE-191:Integer Underflow vulnerability. To priori-
 515 tise the evaluation of the explanation over participant proficiency in detecting
 516 various types of vulnerabilities, only one type of vulnerability was included
 517 in the user study.

518 The hypothesis was evaluated with independent groups of participants
519 recruited through Amazon Mechanical Turk. One group received the ques-
520 tionnaire together with DisCERN counterfactual explanations and the other
521 with LIME feature attribution explanations. From here on we will refer to the
522 two groups as DisCERN and LIME. The inclusion criteria for recruitment
523 were set as *Employment Industry* is *Software and/or IT Services* and *Job*
524 *function* is *Information Technology* to ensure the participants have a working
525 knowledge of programming languages. In 40 days, 95 and 103 submissions
526 were received for DisCERN and LIME groups respectively from which 78
527 and 68 were accepted. These submissions met the minimum requirements
528 where they attempted to answer at least one free-text question in addition
529 to all multiple choice questions (There were only 9 and 12 submissions for
530 DisCERN and LIME groups where participants answered all questions).

531 6.1. *A priori mental model - detecting code vulnerabilities*

532 Q1 measures the a priori mental model for understanding how to detect
533 code vulnerabilities. There are 438 responses (78 + 68 participants responded
534 to 3 code snippets each) considered in total. Figure 7a plots the percentage of
535 *Yes*, *No* and *Maybe* responses from the two groups. The percentages between
536 the groups are comparable which suggests that the a priori knowledge and
537 understanding levels are similar. However, the LIME group demonstrates
538 higher accuracy and more confidence in their decision choices evidenced by
539 the lower percentage in *Maybe* responses.

540 Figure 7b plots the percentage of responses received for each snippet.
541 The DisCERN group identifies Snippet 2 as the most complex, as evidenced
542 by their higher percentage of *Maybe* responses. Additionally, we observe
543 that the high confidence of the LIME group stems from the least complex
544 Snippet 1. Both observations imply that the responses are not arbitrary,
545 lending credibility to the utilisation of Q1 responses as an indicator of the
546 group's a priori mental model.

547 6.2. *A priori mental model - correcting code vulnerabilities*

548 Q2 measures the a priori mental model for correcting code vulnerabili-
549 ties. Participants answered Q2 with line numbers or code lines which they
550 considered to be *vulnerable*. Few example responses were *3,4,5, int data =*
551 *method();* and *3rd line*. After pre-processing, Table 4 plots the number of
552 responses for the three snippets across the two groups against corresponding
553 code lines. Here the number of responses relates to the number of times

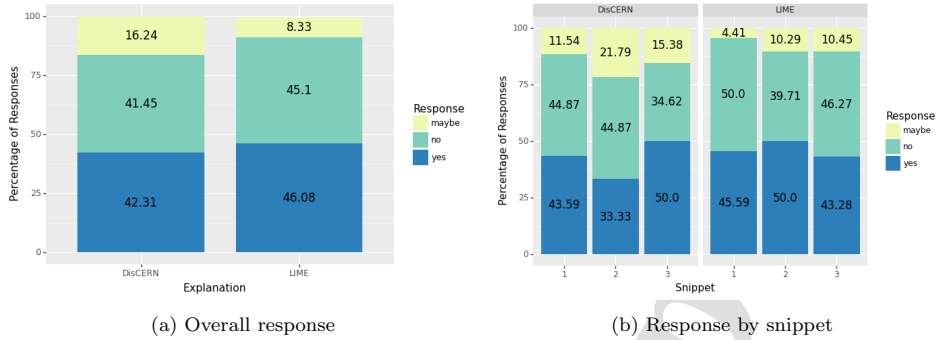


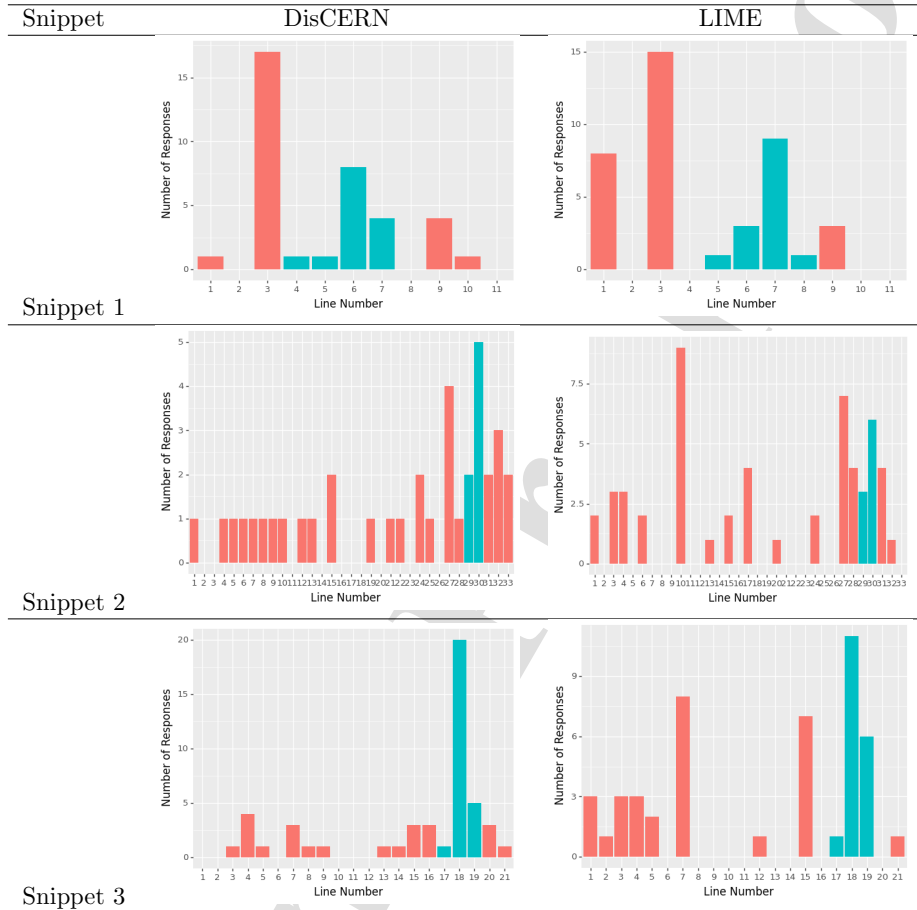
Figure 7: Q1 analysis on a priori mental model - detecting code vulnerabilities

554 a specific line was identified as vulnerable. We then analyse these against
 555 the actual vulnerable lines (the ground truths). The plots use a two-way
 556 colour coding to distinguish between lines that are correctly identified as
 557 vulnerable (in blue) and those that are incorrectly identified as vulnerable
 558 (in red). Although we wouldn't anticipate participants who answered *No* (or
 559 to a lesser extent *Maybe*) in Q1 to respond to Q2, we have still included their
 560 Q2 responses in the graphs if they chose to provide them.

561 We calculate response accuracy as a percentage of correct responses com-
 562 pared to ground truth. DisCERN group demonstrated 37.8%, 18.9%, 53.1%
 563 response accuracy while LIME group achieves 35.0%, 16.7%, 38.3%. Overall
 564 accuracy for DisCERN and LIME groups were 36.6% and 30.0%. Snippet
 565 2 was the most challenging for both groups indicated by the lowest accu-
 566 racy, The wide variety of responses suggests that the increased complexity
 567 made participants uncertain and led to guessing. Overall, guessing or ran-
 568 dom responses are expected from those who did not detect vulnerability in
 569 Q1.

570 We observe that the code segment length has some correlation to the
 571 number of errors. Accordingly, we further normalise the accuracy values by
 572 the “difficulty of predicting vulnerable code lines in a code segment” using
 573 inspirations from document length normalisation which alleviates the “term-
 574 frequency-bias”. Given the number of lines of code in the segment is α out
 575 of which β number of lines are vulnerable, the difficulty is calculated as
 576 $1 - \beta/\alpha$. If all lines were vulnerable $\beta = \alpha$ then *difficulty* = 0 and vice
 577 versa. The weighted accuracy values are 20.4%, 17.8% and 45.6% for Dis-
 578 CERN group (mean is 27.93%) and 18.9%, 15.7% and 32.9% for the LIME

Table 4: Q2 analysis on a priori mental model - correcting code vulnerabilities



579 group (mean is 22.5%). The difference between the two groups is influenced by
 580 two factors: the number of responses for Snippet 2 from the DisCERN group
 581 was significantly lower than LIME group (37 vs 54) which contributed to the
 582 2.1% difference, and for Snippet 3 DisCERN group responses were signifi-
 583 cantly more accurate (45.6% over 32.9%) although the number of responses
 584 was comparable (49 vs 47). This analysis aids in determining the groups'
 585 initial mental models, which is valuable for assessing the subsequent changes
 586 in their mental models a posteriori. We recognise the marginally higher (ap-
 587 proximately 5%) performance of the DisCERN cohort and will consider this

588 in our subsequent analysis when we focus on a posteriori evaluations.

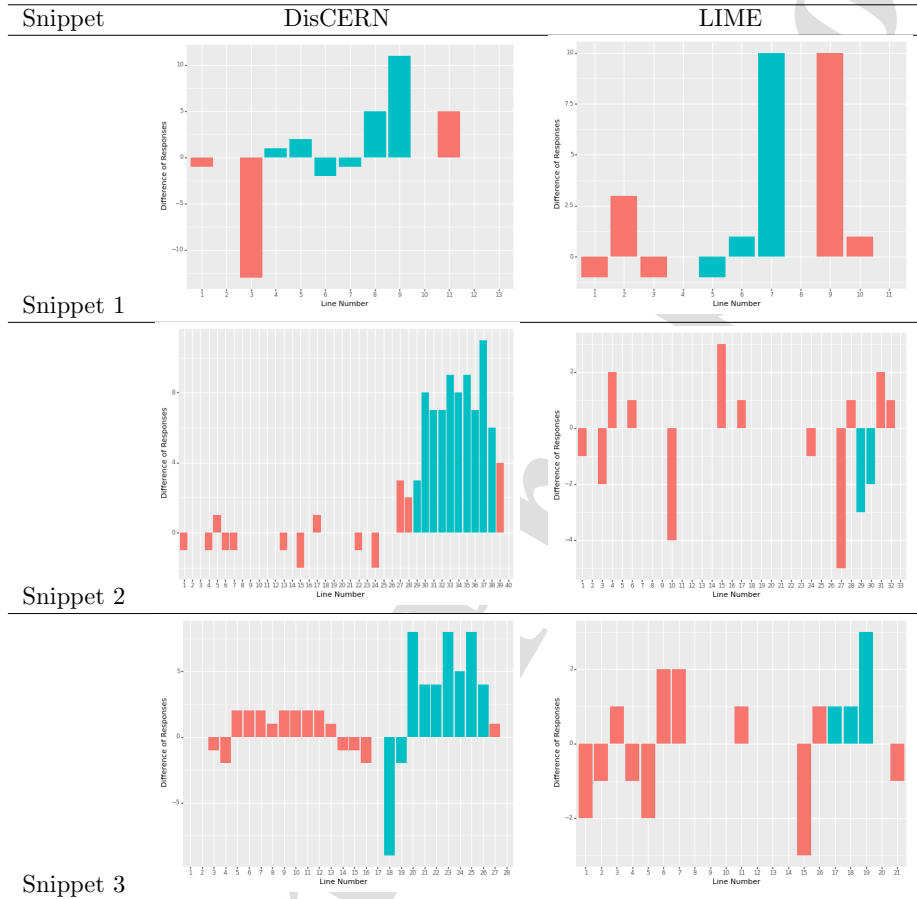
589 *6.3. A posteriori mental model for correcting code vulnerabilities*

590 Q4 measures the a posteriori mental model for addressing vulnerabilities
591 after participants have been exposed to the explanation. This implies that
592 participants have been informed about the snippet’s vulnerability and are
593 presented with an explanation—either a counterfactual from DisCERN or
594 feature attribution from LIME. The explanations were presented as code-diff
595 for DisCERN and heat maps for LIME. To minimise the possibility of misin-
596 terpretations, we have provided supporting text alongside both explanations,
597 detailing how to interpret them effectively.

598 Following pre-processing of the participants’ responses, we analysed any
599 changes in knowledge among each user group after exposure to the explana-
600 tion for the three code snippets, as shown in Table 5. Here, we anticipate
601 that changes to their mental model will be evident in at least two ways: 1)
602 withdrawing their belief for lines that were incorrectly identified as vulner-
603 able in Q2, and 2) recognising new lines that are necessary to address the
604 vulnerability having seen the explanation in Q4. For example, if the change
605 in response for a code line is denoted by -3 , it means that the number of
606 responses for that line after participants saw the explanation (Q4) decreased
607 by three compared to before (Q2), indicating a shift in their belief about
608 the vulnerability of that line. Here, the reductions observed with the Orange
609 lines represent a positive change that was achieved a posteriori. Unlike LIME,
610 DisCERN not only identifies vulnerabilities but also provides hints on how
611 to correct them by displaying counterfactuals. As a result, participants can
612 access additional lines from the counterfactual that were not available in Q2.
613 This is seen in Table 5 for DisCERN, where a relatively larger number of
614 blue lines can be observed on the x-axis, indicating a notable difference.

615 Overall Table 5 observations strongly indicate that participants found
616 counterfactuals more informative to correct vulnerabilities compared to fea-
617 ture attributions. The DisCERN group exhibited some errors, as misiden-
618 tified lines on either side of the vulnerability boundary were observed. For
619 instance, in Snippet 2, lines 28 and 39 were considered worthy of change,
620 despite not being vulnerable. Similarly, in Snippet 3, lines 18 and 19 were
621 not recognised as vulnerable, representing another error. The boundary cases
622 observed with DisCERN and the errors observed with the LIME group both
623 suggest that some participants are likely to either misinterpret or disagree
624 with the explanations.

Table 5: Q4 analysis on a posterior mental model - correcting code vulnerabilities



625 6.4. Goodness of explanations for vulnerability detection and correction

626 Q6 and Q7 aim to measure the overall goodness of the explanation to
 627 detect and correct vulnerabilities. Both questions are further analysed in
 628 relation to Q1 to examine the utility of the explanations to different cohorts:
 629 knowledgeable participants who responded *Yes* in Q1; and trainee partici-
 630 pants who responded *No* or *Maybe* in Q1.

631 Q6 results across the two groups are plotted in Figure 8. The posi-
 632 tive response rate from DisCERN and LIME groups were 66.7% and 62.7%
 633 respectively when asked about the utility of explanations for vulnerability

634 detection. This indicates a slight preference towards counterfactual explanations.
 635 Furthermore, Figure 8b indicates that the counterfactual explanations
 636 were found to be useful for more complex snippets (2 and 3) and feature
 637 attributions useful for the smallest snippet (1). This suggests that using
 638 counterfactual explanations may result in a lower cognitive load for detect-
 639 ing errors when compared to feature attributions.

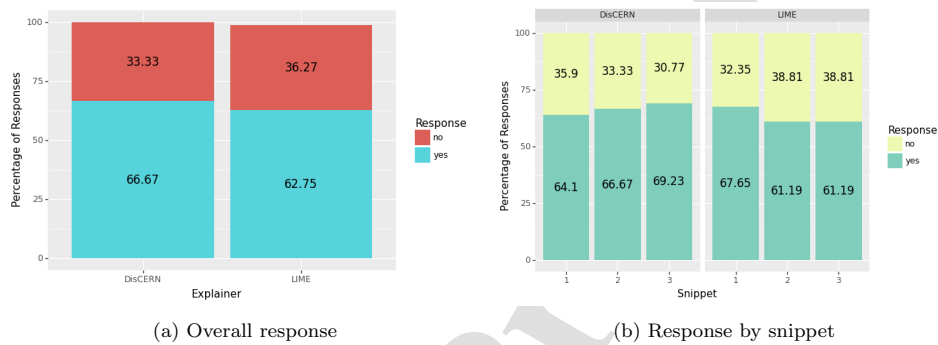


Figure 8: Q6 analysis on the goodness of explanations - detecting code vulnerabilities

640 Figures 9a and 9b present an in-depth analysis of the Q6 responses with
 641 respect to Q1. Figure 9a shows that participants with prior knowledge of
 642 vulnerability detection found both types of explanations useful. The im-
 643 proved positive response rates of 83.61% and 77.56% from their baselines
 644 for DisCERN and LIME indicate that knowledgeable users found both types
 645 of explanations helpful. However, counterfactuals have been significantly
 646 more helpful than feature attributions, especially for complex code snippets.
 647 Figure 9b shows that trainee cohorts struggle with types of explanations.
 648 It is indicated by the decreased positive response rate from their baselines
 649 to 53.7% and 50.2% for DisCERN and LIME groups. However, trainee co-
 650 horts found counterfactuals significantly helpful for the most complex snippet
 651 whereas feature attribution helped with the simplest snippet. These obser-
 652 vations further verify that counterfactuals reduced the cognitive burden of
 653 vulnerability detection in complex code snippets.

654 Q7 measures the utility of the explanation to **correct** vulnerabilities and
 655 we plot similar graphs to Q6. Figure 10a shows that the overall positive
 656 response rates from DisCERN and LIME groups were 66.24% and 63.24%
 657 respectively. Similar to detection (Q6), the responses for Q7 indicate a pre-
 658 ference for the counterfactuals for more complex snippets. In contrast to

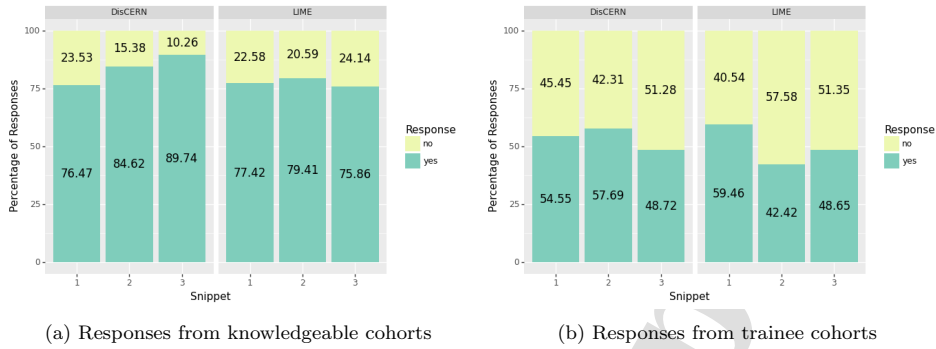


Figure 9: Q6 analysis on the goodness of explanations - detecting code vulnerabilities by knowledgeable and trainee cohorts

659 detection, counterfactuals are found to be comparably helpful for the correc-
 660 tion of vulnerabilities in simpler snippets which evidence an overall preference
 661 towards counterfactual explanations.

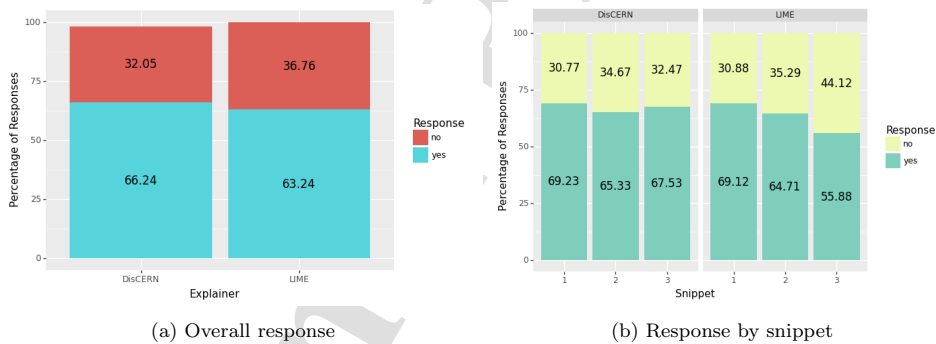


Figure 10: Q7 analysis on the goodness of Explanations - correcting code vulnerabilities

662 Figure 11 presents the analysis of the Q7 response with respect to cohorts
 663 identified in Q1. Similar to Q6, the positive response rate for both explana-
 664 tions have improved from the knowledgeable cohort and decreased from the
 665 trainee cohort. The preference for counterfactuals over feature attributions
 666 by both cohorts for complex snippets remains significant for vulnerability correc-
 667 tion. While trainee cohorts consistently find counterfactuals to be more
 668 helpful, knowledgeable cohorts find feature attributions are sufficient for cor-
 669 recting vulnerabilities in simple snippets.

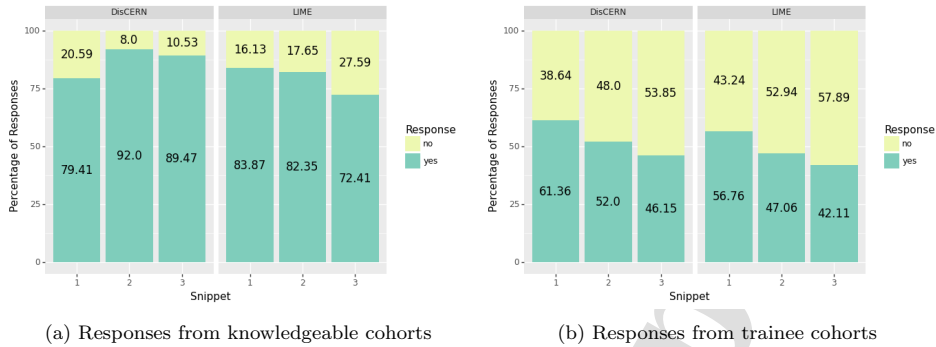


Figure 11: Q7 analysis on the goodness of explanations - correcting code vulnerabilities by knowledgeable and trainee cohorts

670 We acknowledge that a significant number of trainee users did not find
 671 either type of explanation helpful for both detection and correction. This was
 672 clearly seen from responses to both Q6 and Q7 having less than 50% positive
 673 responses for Snippets 2,3 in the LIME group and Snippet 3 in DisCERN
 674 group. However, the overall preference was for DisCERN counterfactuals.
 675 This feedback is useful for future research to further improve counterfactual
 676 explanations to assist trainee developers to learn about vulnerability detec-
 677 tion and correction.

678 6.5. Acceptability of the explanations

679 Q8 is aimed to measure the acceptability of the explanations. Similar to
 680 Q6 and Q7 we plot overall responses and responses by snippets in Figure 12.
 681 In both groups, approximately 60% of the participants agreed with the ex-
 682 planations provided. However, the disagreement is significantly lower in the
 683 DisCERN group where 3% more partially accepted the counterfactual ex-
 684 planation. Figure 12b shows that the acceptability of LIME is significantly
 685 lower for Snippet 3 which has affected the overall acceptance. Otherwise,
 686 agreement with feature attributions is similar to or greater than that of coun-
 687 terfactuals which is inconsistent with the previous observations on change in
 688 mental model and goodness. LIME is a well-established explanation method
 689 in various domains for several years, which may have influenced the observed
 690 results, to further verify, we perform a more in-depth analysis.

691 The first analysis of acceptability is with respect to the cohorts recognised
 692 in Q1. Figure 13 plots the acceptability by knowledgeable and trainee co-

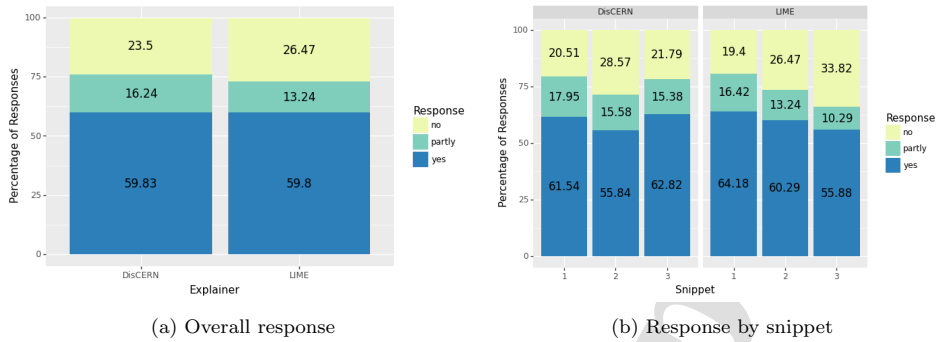


Figure 12: Q8 analysis on the acceptability of explanations

693 horts. The knowledgeable cohorts found counterfactuals more agreeable than
 694 feature attributions, indicated by the accumulated positive response rates of
 695 81.47% and 75.60%. The most significant difference is that the counterfactual
 696 explanation for the most complex snippet is found to be more agreeable
 697 which aligns with previous observations. We failed to observe a majority of
 698 trainee cohorts agreeing with either explanation, however, we observe partial
 699 agreement rates of 63.70% and 62.20% respectively for DisCERN and LIME.
 700 These findings reinforce the overall utility of counterfactuals over feature at-
 701 tribution and also highlight the need to improve the counterfactuals to build
 702 trust among trainee developers as an effective learning tool.

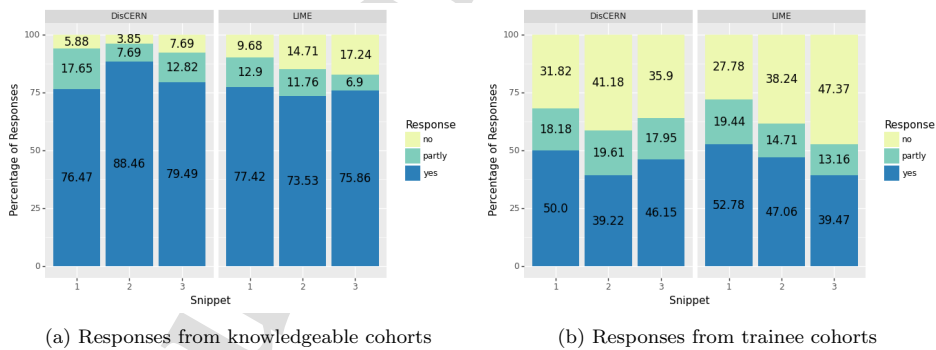
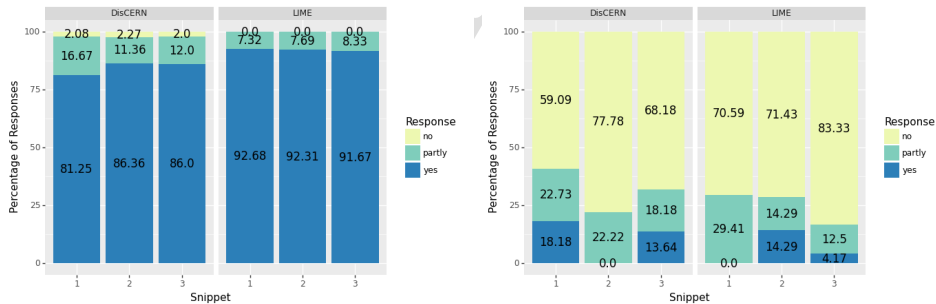


Figure 13: Q8 analysis on the acceptability of explanations by knowledgeable and trainee cohorts

703 The second analysis of acceptability is with respect to the explanation
 704 goodness observed by Q6 and Q7. We recognise two cohorts from Q6 and

705 Q7, the ones who found explanations helpful and others who did not for both
 706 vulnerability detection and correction. Figure 14a plots the Q8 responses for
 707 those who found explanations helpful. Results show that the participants
 708 who found feature attribution helpful overwhelmingly agreed with the expla-
 709 nation (0% *no* responses). However, not all who found counterfactuals use-
 710 ful agreed with it indicated by 2.11% disagreeing and 13.34% only partially
 711 agreeing. Figure 14b plots the Q8 responses for those who found explana-
 712 tions not helpful. Those who found feature attribution not helpful for small
 713 snippets completely disagreed with the explanation and those who found
 714 counterfactuals not helpful for complex snippets, also completely disagreed
 715 with the explanation. It is noteworthy that both of these cohorts are the
 716 minority when determining goodness. Overall, 31.63% and 24.89% at least
 717 partially agreed with counterfactuals and feature attributions respectively.

718 These observations conclude that the higher overall agreement with fea-
 719 ture attributions seen in Figure 12b for Snippets 1 and 2 is influenced by
 720 the cohorts who found counterfactuals helpful (Q6 and Q7) but did not fully
 721 agree with them. What is unknown and needs to be established in the long
 722 term is if this acceptance of feature attribution is influenced by familiarity
 723 with LIME explanations. The need for this is supported by the results in
 724 Section 6.3 which clearly showed that counterfactuals influence a positive
 725 mental model change compared to feature attributions.



(a) Responses from Q6, Q7 = yes cohorts

(b) Responses from Q6, Q7 = no cohorts

Figure 14: Q8 analysis on the acceptability of explanations by goodness cohorts

726 6.6. Implications and Limitations of the User Study

727 Overall, counterfactual explanations encourage positive mental model
 728 changes and were perceived as more helpful than feature attributions for

729 detecting and correcting code vulnerabilities. However, feature attributions
730 exhibit comparable or higher acceptability, possibly due to their widespread
731 use. These conclusions should be made with the limitations of the study in
732 mind. The key limitations of the above user study are three-fold: 1) incom-
733 pleteness of heuristics used to identify the knowledgeable and trainee cohorts;
734 2) inclusion and exclusion criteria of participants; and 3) representations and
735 interpretation of explanations.

736 The a priori mental model for detecting vulnerabilities was based on Q1
737 which recognised two cohorts as knowledgeable and trainee. However, we did
738 not account for those who recognised the vulnerabilities incorrectly by col-
739 lating them with answers to Q2. As seen in Section 6.2 only 36% and 30% of
740 the two groups identified the lines correctly and it includes all participants.
741 We found it challenging to filter participants by both Q1 and Q2 because
742 many of those who answered *Yes* in Q1 were able to partially identify vul-
743 nerable lines. A more strict filter would have resulted in no knowledgeable
744 participants. Accordingly, we relied solely on Q1 to categorise participants
745 into the two cohorts.

746 The recruitment of participants for the user study was limited to Amazon
747 Mechanical Turk (AMT), which placed constraints on the inclusion criteria.
748 Accordingly, the inclusion criteria for selecting participants were constrained
749 to those possible in the AMT platform. Ideally, a more comprehensive study
750 would include participants in various career stages with Java software devel-
751 opment skills.

752 The explanations generated by LIME and DisCERN are significantly dif-
753 ferent in their presentation. LIME highlights the original query using a
754 heat-map scale and DisCERN presents counterfactual in a code-diff view.
755 With LIME explanations where the individual tokens are highlighted, it may
756 mislead the participants. An example scenario is if two tokens in a state-
757 ment were highlighted as *vulnerable* and *non-vulnerable*, the participant can
758 consider the statement as *vulnerable*, *non-vulnerable* or *have no impact* on
759 the vulnerability detection. DisCERN code-diff view has two columns with
760 query line numbers and counterfactual line numbers. A participant who is
761 unfamiliar with code-diff may mistakenly use the line numbers from the in-
762 appropriate column when responding to the questionnaire.

763 All three limitations are well-founded, however, they do not invalidate
764 the findings, rather, they provide enhancing user studies in this particular
765 domain and in Explainable AI in general.

766 The user study was conducted with three code segments, all of which be-

767 longed to the same vulnerability code (CWE-191). Our main reasoning was
768 to prioritise the evaluation of the utility of different types of explanation while
769 keeping other variables constant. Additionally, it allowed the user study not
770 to be biased by the proficiency of the participant in detecting various types of
771 vulnerabilities. However, there can be implications for this approach if some
772 vulnerabilities were better explained using feature attributions over counter-
773 factuals. This can be linked to our observations in Section 6.4 where there
774 was no significant preference between feature attribution and counterfactual
775 explanations when the code segment was simple. The generalisability of Dis-
776 CERN to different vulnerability classes and languages (as seen in Section 5)
777 provides an opportunity to evaluate this in the future.

778 7. Conclusion

779 The DisCERN algorithm finds counterfactual explanations for correcting
780 code vulnerabilities using pattern matching to find corrections to a code seg-
781 ment from its nearest-unlike neighbour. DisCERN was evaluated using three
782 NIST datasets in different programming languages and the results showed
783 that it finds counterfactuals in 85% - 96% of the cases with 8 ~ 14 statement
784 corrections needed. A qualitative analysis revealed that some of the counter-
785 factuals generated by DisCERN did not preserve the original functionality of
786 the code. This highlights the need for comprehensive heuristics in the future
787 to ensure plausible code corrections. We conducted a user study to assess the
788 utility of counterfactual explanations compared to the more commonly used
789 feature attribution explanations for correcting vulnerabilities. The user study
790 showed that counterfactuals facilitated a positive mental model change to-
791 wards correcting vulnerabilities. Counterfactuals were specifically preferred
792 over feature attributions when dealing with complex code segments, indi-
793 cating a reduction in cognitive burden. However, despite being less helpful
794 for vulnerability correction, feature attribution explanations received higher
795 acceptance than counterfactuals, possibly due to the trust built around their
796 familiarity. These findings provide evidence for the utility of counterfactual
797 explanations over feature attribution explanations. Nonetheless, they also
798 emphasise the importance of conducting long-term evaluations to determine
799 if counterfactuals can establish trust with developers as a reliable tool for
800 vulnerability detection and correction.

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 803 project ⁶ which received funding for the UK from EPSRC under the grant
 804 number EP/V061755/1.

805 **Appendix A. Code snippets from the user study**

```
public void method()
{
  int data = someMethod();

  /*comment*/
  data--;
  int result = (int)(data);

  IO.writeLine("result: " + result);
}
```

Figure A.15: Snippet 1: LIME Explanation

```
1 1 public void method()
2 2 {
3 -   int data = someMethod();
3+  int data;
4 4
5 5   /*comment*/
6 -   data--;
7 -   int result = (int)(data);
6+  data = 2;
7+
8+   /*comment*/
9+   int result = (int)(--data);
8 10
9 11   IO.writeLine("result: " + result);
10 12
11 13 }
```

Figure A.16: Snippet 1: DisCERN Explanation

⁵<https://isee4xai.com/>

⁶<https://www.chistera.eu/projects/isee>

```
public void method()
{
    int data;
    if (IO.STATIC_FINAL_TRUE)
    {
        data = Integer.MIN_VALUE; /*comment*/
        /*comment*/
        /*comment*/
        {
            String stringNumber = System.getProperty("user.home");
            try
            {
                data = Integer.parseInt(stringNumber.trim());
            }
            catch(NumberFormatException exceptNumberFormat)
            {
                IO.logger.log(Level.WARNING, "Number format exception parsing data from string",
                exceptNumberFormat);
            }
        }
    }
    else
    {
        /*comment*/
        data = 0;
    }
    if (IO.STATIC_FINAL_TRUE)
    {
        /*comment*/
        int result = (int)(data - 1);
        IO.writeLine("result: " + result);
    }
}
```

Figure A.17: Snippet 2: LIME Explanation

```
1 1 public void method()
2 2 {
3 3     int data;
4 4     if (IO.STATIC_FINAL_TRUE)
5 5     {
6 6         data = Integer.MIN_VALUE; /*comment*/
7 7         /*comment*/
8 8         /*comment*/
9 9         {
10 10            String stringNumber = System.getProperty("user.home");
11 11            try
12 12            {
13 13                data = Integer.parseInt(stringNumber.trim());
14 14            }
15 15            catch(NumberFormatException exceptNumberFormat)
16 16            {
17 17                IO.logger.log(Level.WARNING, "Number format exception parsing data from string", exceptNumberFormat);
18 18            }
19 19        }
20 20    }
21 21    else
22 22    {
23 23        /*comment*/
24 24        data = 0;
25 25    }
26 26
27 27    if (IO.STATIC_FINAL_TRUE)
28 28    {
29 29        | /*comment*/
30 -        int result = (int)(data - 1);
31 -        IO.WriteLine("result: " + result);
30+        if (data > Integer.MIN_VALUE)
31+        {
32+            int result = (int)(data - 1);
33+            IO.WriteLine("result: " + result);
34+        }
35+        else
36+        {
37+            IO.WriteLine("");
38+        }
32 39    }
33 40 }
```

Figure A.18: Snippet 2: DisCERN Explanation

```

public void method()
{
    byte data;
    if (IO.staticTrue)
    {
        /*comment*/
        data = Byte.MIN_VALUE;
    }
    else
    {
        /*comment*/
        data = 0;
    }

    if (IO.staticTrue)
    {
        /*comment*/
        byte result = (byte)(data - 1);
        IO.writeLine("result: " + result);
    }
}

```

Figure A.19: Snippet 3: LIME Explanation

```

1 1 public void method()
2 2 {
3 3     byte data;
4 4     if (IO.staticTrue)
5 5     {
6 6         /*comment*/
7 7         data = Byte.MIN_VALUE;
8 8     }
9 9     else
10 10    {
11 11        /*comment*/
12 12        data = 0;
13 13    }
14 14
15 15    if (IO.staticTrue)
16 16    {
17 17        /*comment*/
18 -        byte result = (byte)(data - 1);
19 -        IO.writeLine("result: " + result);
18+        if (data > Byte.MIN_VALUE)
19+        {
20+            byte result = (byte)(data - 1);
21+            IO.writeLine("result: " + result);
22+        }
23+        else
24+        {
25+            IO.writeLine("");
26+        }
20 27    }
21 28 }

```

Figure A.20: Snippet 3: DisCERN Explanation

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