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**Decision problems concerning a power series generalization of DT0L systems. (English. English summary)**

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Let  $\Sigma$  be an alphabet,  $A$  a commutative semiring, and  $A\langle\Sigma^*\rangle$  the  $A$ -semialgebra of formal polynomials with noncommuting variables from  $\Sigma$  and coefficients from  $A$ . A 3-tuple  $G = (\Sigma, H, v)$  is a “DT0L system with multiplicities in  $A$ ” if  $v \in A\langle\Sigma^*\rangle$  is a nonzero monomial and  $H = \{h_1, \dots, h_n\}$ , where  $h_1, \dots, h_n: A\langle\Sigma^*\rangle \rightarrow A\langle\Sigma^*\rangle$  are monomial morphisms (i.e., for each  $\sigma \in \Sigma$  there exist a nonzero  $a \in A$  and a  $w \in \Sigma^*$  such that  $h_i(\sigma) = aw$  for  $1 \leq i \leq n$ ). The mapping  $S(G): H^* \rightarrow A\langle\Sigma^*\rangle$ , defined by  $S(G)(u) = u(v)$ , is the “DT0L sequence with multiplicities generated by  $G$ ”.

The author establishes (i) a general result (that each recursively enumerable set over  $\mathbf{N}$  can be characterized via a multiplicity sequence defined by a DT0L system with multiplicities in  $\mathbf{Q}$ ), which suggests that most questions concerning DT0L power series and multiplicity sequences are undecidable, and—to the contrary—(ii) the sequence equivalence problem (i.e., “Is  $S(G_1) = S(G_2)$ ?”) is decidable for DT0L systems with multiplicities in  $\mathbf{Q}$ . *Peter R. J. Asveld* (NL-TWEN-C)