

Characterizing Coherence, Correcting Incoherence



1. Context

Basic setup:

- Finite possibility space Ω
- Finite set of gambles \mathcal{K} on Ω
- Lower previsions \underline{P} on \mathcal{K}

Matrix notation:

- $|\Omega|$ -by- $|\mathcal{K}|$ matrix K with gambles as columns
- the rows of K (columns of K^T) are the degenerate previsions
- the set \mathcal{S} of matrices S obtained from the identity matrix \mathbb{I} by changing at most one 1 to -1
- all-one (zero) column vector $\mathbf{1}$ ($\mathbf{0}$)

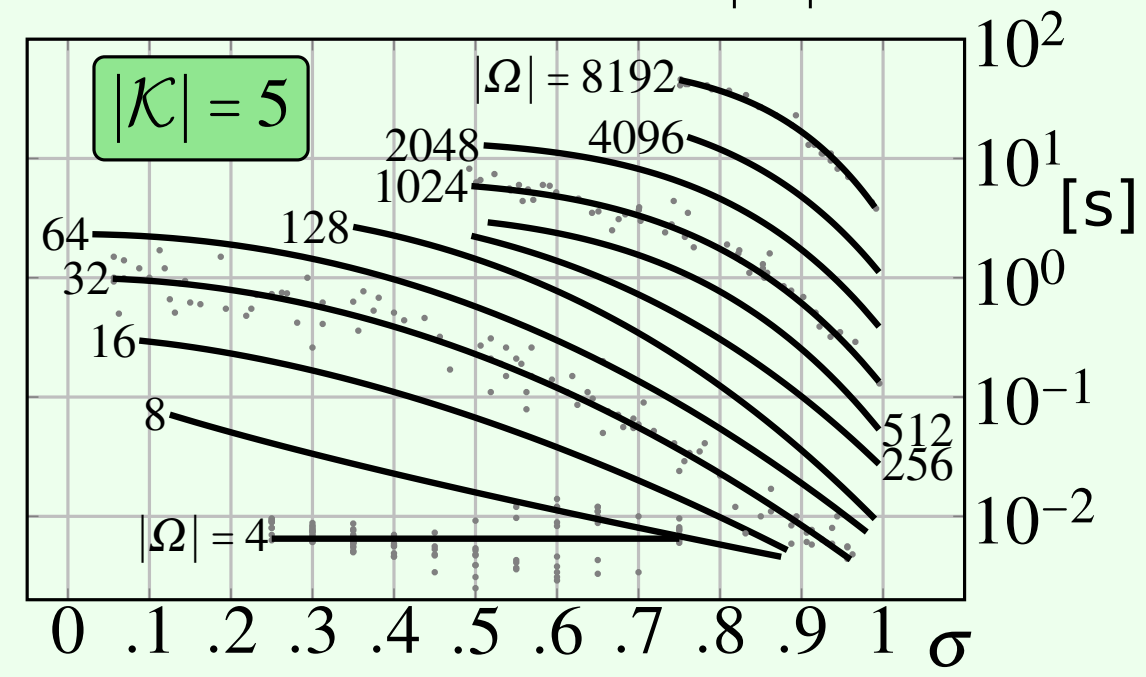
2. Goals

- Given K , find a non-redundant H-representations for the set of all \underline{P}
- A. that avoid sure loss ($[\Lambda_A | \alpha_A]$),
 - B. that avoid sure loss and for which $\underline{P} \geq \min$ ($[\Lambda_B | \alpha_B]$),
 - C. that are coherent ($[\Lambda_C | \alpha_C]$).

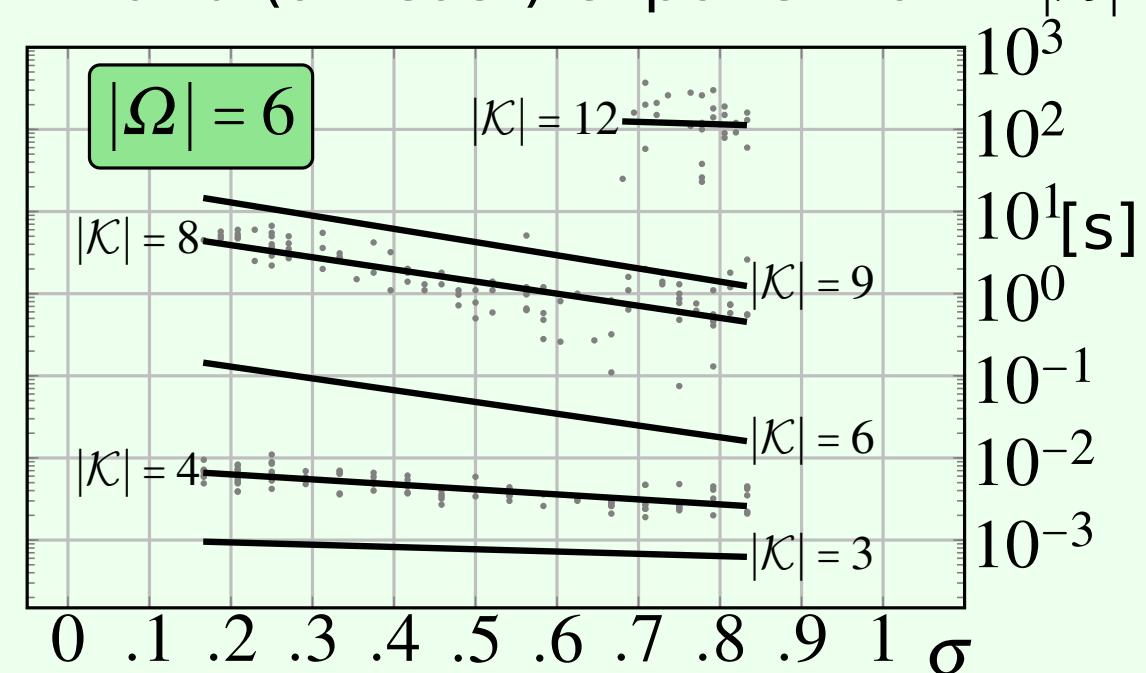
7. Experiments

The sparsity σ is the fraction of zero components in K .

Procedure C1 is exponential in $1 - \sigma$ and \sim linear in $|\Omega|$:



... and (at least) exponential in $|\mathcal{K}|$:



3. Goal A: Characterizing ASL

Based on the existence of a dominating linear prevision:

$$A1. \exists \mu_{\mathbb{I}}, \nu_{\mathbb{I}} \geq 0: \underline{P} = K^T \mu_{\mathbb{I}} - \mathbb{I} \nu_{\mathbb{I}} \wedge \mathbf{1}^T \mu_{\mathbb{I}} = 1 \xrightarrow{EN, RR} [\Lambda_A | \alpha_A]$$

$$A2. \exists \mu_{\mathbb{I}} \geq 0: \underline{P} \leq K^T \mu_{\mathbb{I}} \wedge \mathbf{1}^T \mu_{\mathbb{I}} = 1 \xrightarrow{PJ_P, RR} [\Lambda_A | \alpha_A]$$

4. Goal B: Characterizing ASL $\geq \min$

$$B1. \text{Starting from } [\Lambda_A | \alpha_A]: \begin{bmatrix} \Lambda_A & \alpha_A \\ -\mathbb{I} & -\min \end{bmatrix} \xrightarrow{RR} [\Lambda_B | \alpha_B]$$

5. Goal C: Characterizing coherence

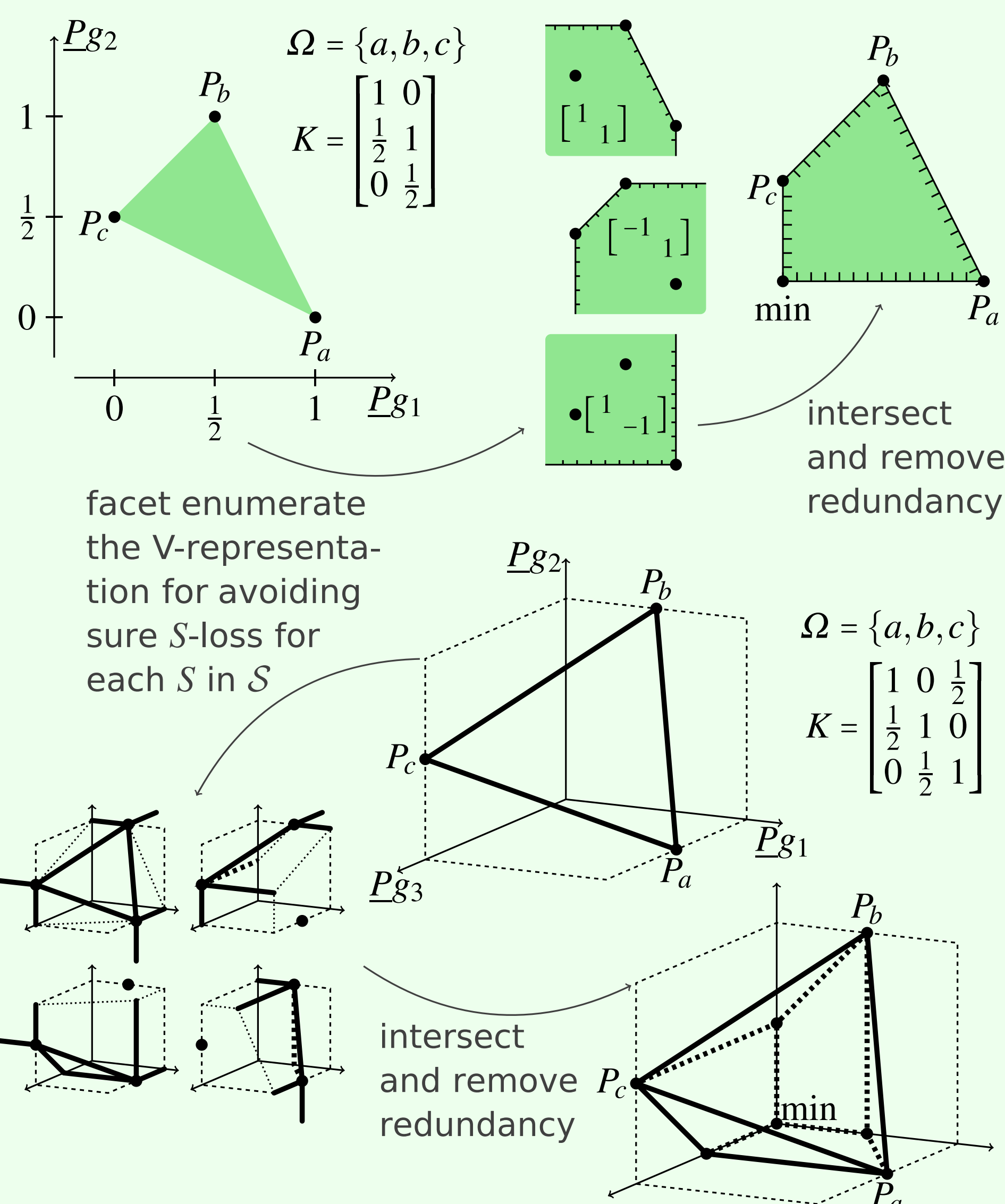
Based on the existence of S -dominating linear previsions:

$$C1. \text{Analogous to A1 \& intersection over all } S \text{ in } \mathcal{S}: \forall S \in \mathcal{S}: \exists \mu_S, \nu_S \geq 0: \underline{P} = K^T \mu_S - S \nu_S \wedge \mathbf{1}^T \mu_S = 1 \xrightarrow{EN, IS_{S \in \mathcal{S}}, RR} [\Lambda_C | \alpha_C]$$

$$C2. \text{Analogous to A2 \& intersection over all } S \text{ in } \mathcal{S}: \forall S \in \mathcal{S}: \exists \mu_S \geq 0: S \underline{P} \leq S K^T \mu_S \wedge \mathbf{1}^T \mu_S = 1 \xrightarrow{PJ_P, IS_{S \in \mathcal{S}}, RR} [\Lambda_C | \alpha_C]$$

$$C3. \text{Block matrix form of C2: } [A_{\underline{P}} \ A_{\mu} \ | \ b] := \begin{bmatrix} A_{\mathbb{I}, \underline{P}} & A_{\mathbb{I}, \mu_{\mathbb{I}}} & b_0 \\ \vdots & \vdots & \vdots \\ A_{S, \underline{P}} & A_{S, \mu_S} & b_0 \\ \vdots & \vdots & \vdots \end{bmatrix} \xrightarrow{PJ_P, RR} [\Lambda_C | \alpha_C]$$

6. Illustrations of Procedure C1



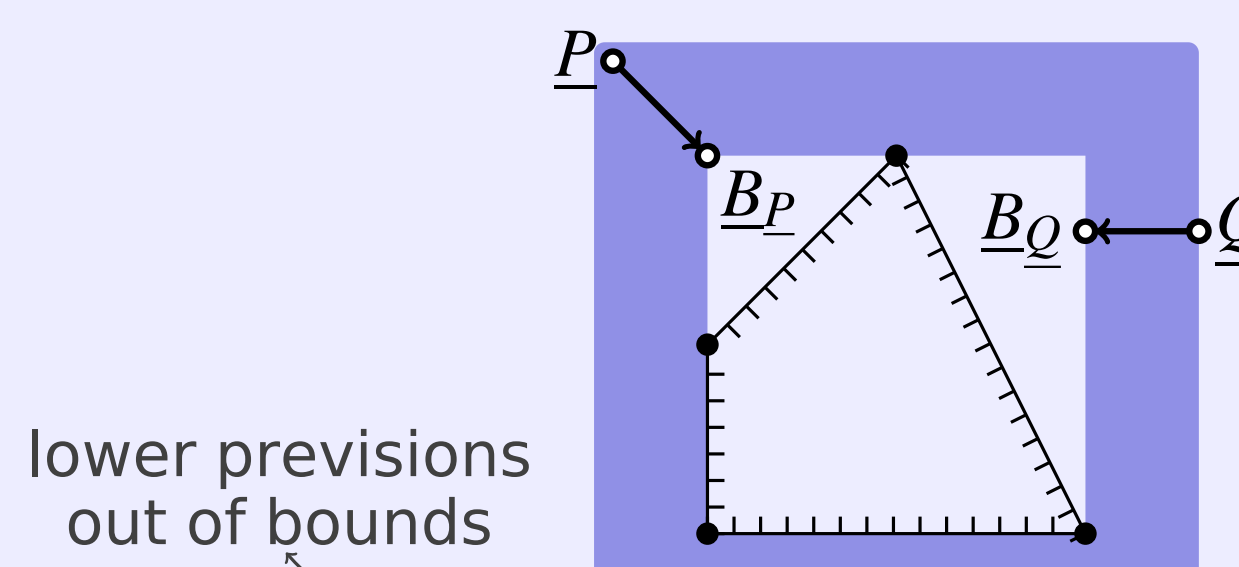
1. Context & Goal

Given: incoherent lower prevision \underline{P} .
Goal: Find a coherent correction to it.

2. Bring within bounds

If $\underline{P}f \notin [\min f, \max f]$ for some f in \mathcal{K} , it is out of bounds. To bring it within bounds:

$$\underline{B}_P f := \begin{cases} \min f & \underline{P}f \leq \min f, \\ \max f & \underline{P}f \geq \max f, \\ \underline{P}f & \text{otherwise.} \end{cases}$$



3. Downward correction

As the downward correction of \underline{P} we take the lower envelope of the maximal coherent dominated lower previsions (proposed earlier by Pelessoni & Vicig, following Weichselberger), so the nadir point \underline{D}_P of the MOLP (cf. C)

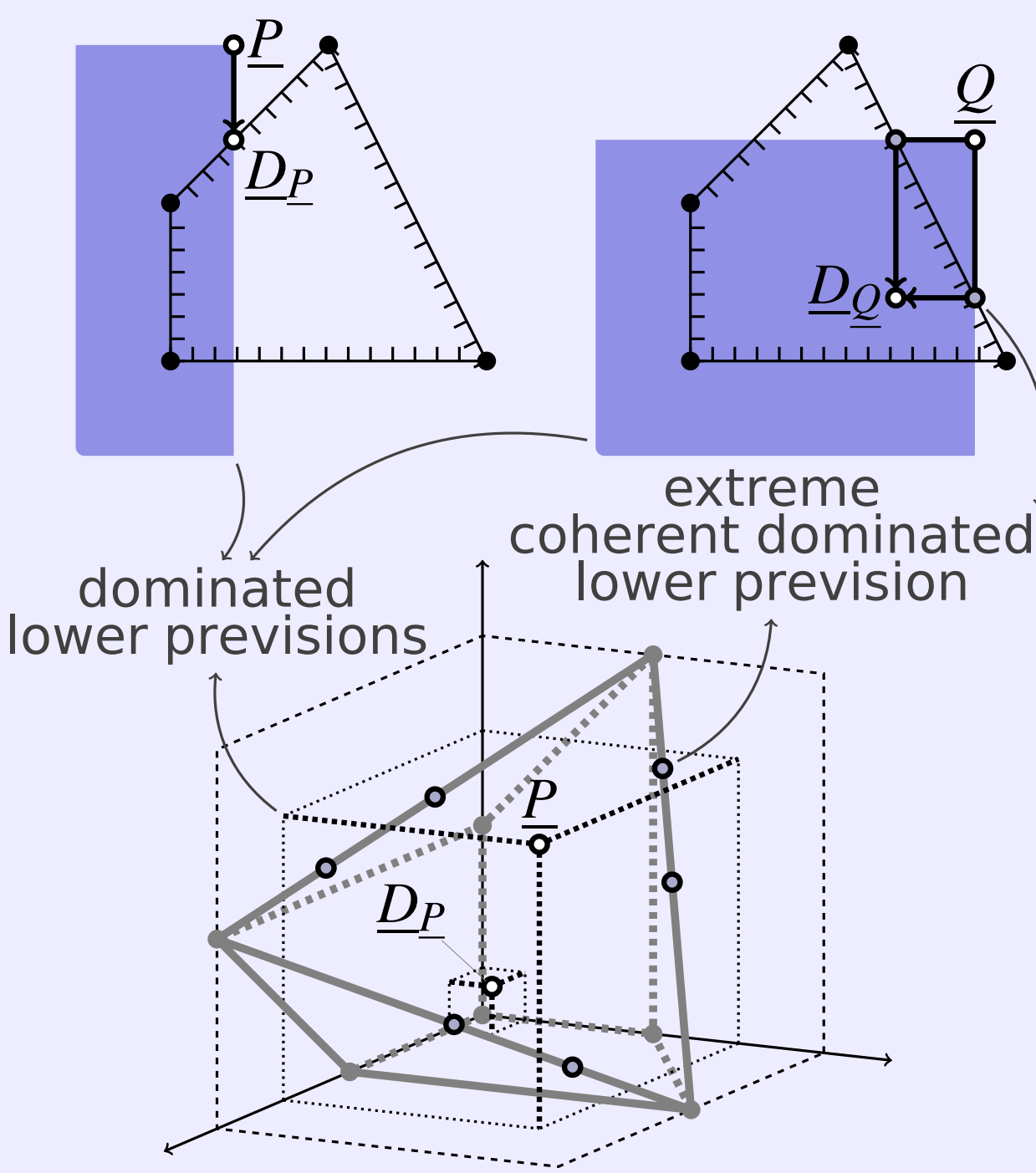
$$\begin{aligned} & \text{maximize } \underline{Q}, \\ (\dagger) \quad & \text{subject to } \Lambda_C \underline{Q} \leq \alpha_C \\ & \underline{Q} \leq \underline{P} \end{aligned}$$

or the MOLP (cf. C3)

$$\begin{aligned} & \text{maximize } \underline{Q}, \\ (\ddagger) \quad & \text{subject to } A_Q \underline{Q} + A_{\mu} \mu \leq b \\ & \underline{Q} \leq \underline{P}. \end{aligned}$$

Some desirable properties:

- It is the maximal neutral correction ('no component tradeoffs').
- The imposition of the correction is nondecreasing with incoherence.

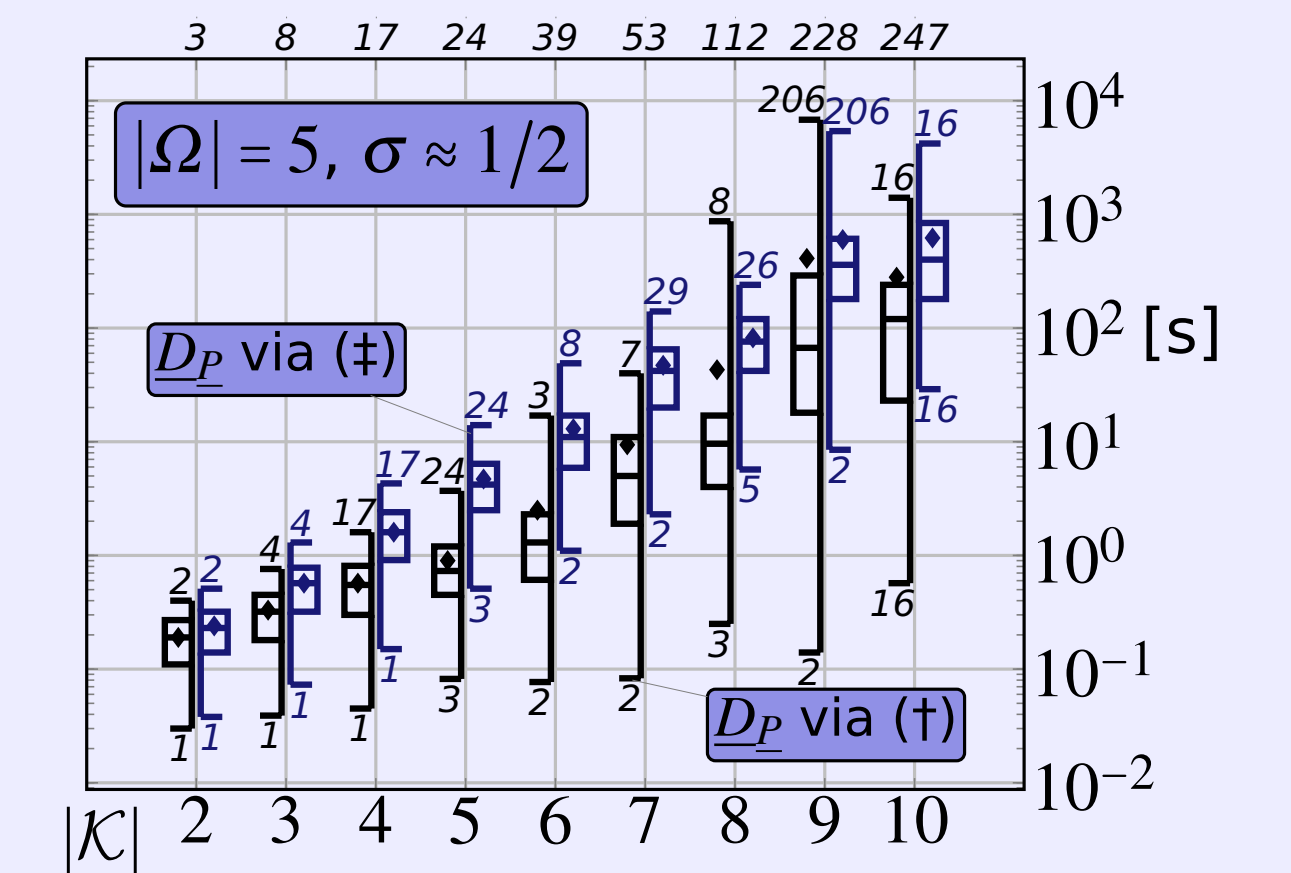


For the future: Can the computation be simplified for special classes of \underline{P} ?



4. Experiments

With the M3-solver we used, computation appears exponential in $|\mathcal{K}|$; using pre-computed constraints (\ddagger) is more efficient than not (\ddagger):



We expect other solvers and certainly direct M2-solvers to perform more efficiently, but could not test any yet.

5. Upward correction

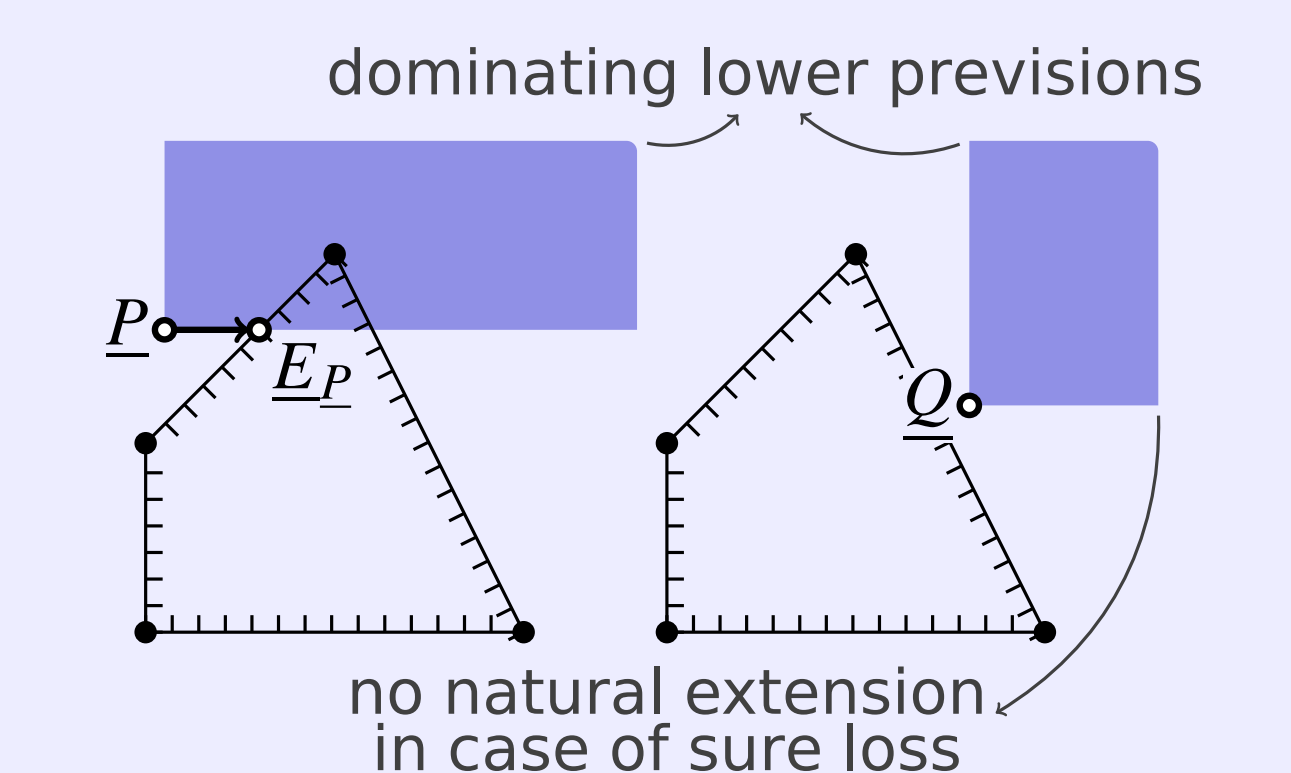
The standard upward correction of \underline{P} is its natural extension \underline{E}_P , the unique minimal pointwise dominating coherent lower prevision, so the the solution to the MOLP (cf. C)

$$\begin{aligned} & \text{minimize } \underline{E}_P, \\ & \text{subject to } \Lambda_C \underline{E}_P \leq \alpha_C \\ & \underline{E}_P \geq \underline{P} \end{aligned}$$

or the MOLP (cf. C3)

$$\begin{aligned} & \text{minimize } \underline{E}_P, \\ (*) \quad & \text{subject to } A_{E_P} \underline{E}_P + A_{\mu} \mu \leq b \\ & \underline{E}_P \geq \underline{P}. \end{aligned}$$

- The problem becomes a plain LP by using the objective $\sum_{g \in \mathcal{K}} E_P g$.
- (*) decomposes into a classical formulation of natural extension.



1. Representations

Any convex polyhedron in \mathbb{R}^n can be described in two ways:

H-representation (intersection of half-spaces)

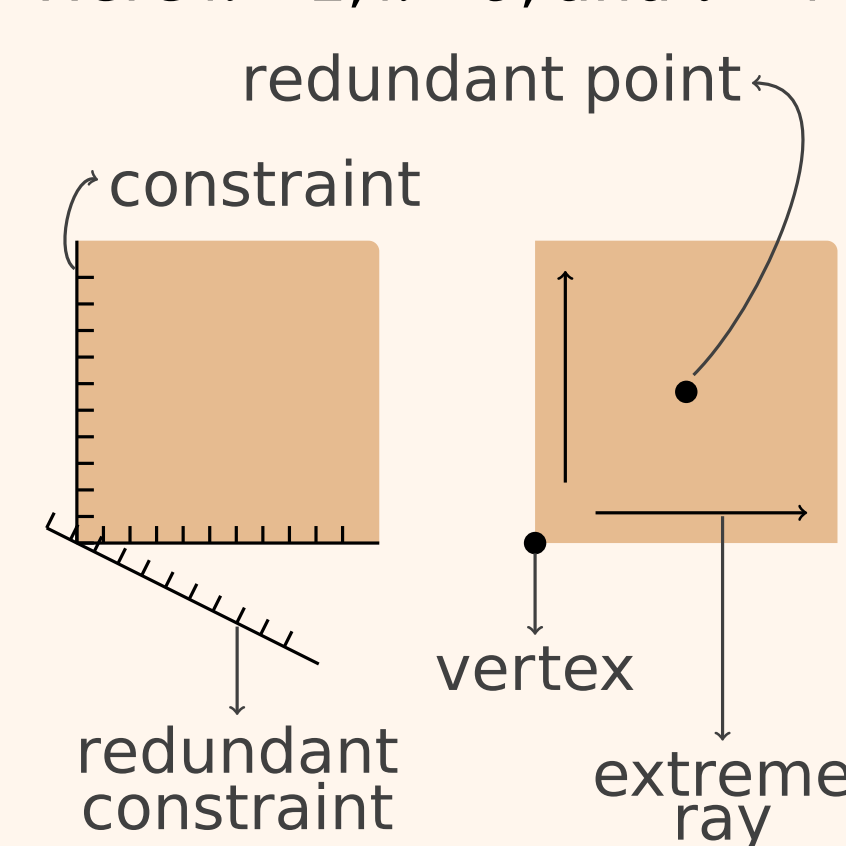
$$[A | b] := \{x \in \mathbb{R}^n : Ax \leq b\}$$

V-representation (convex hull of points and rays)

$$\begin{bmatrix} V \\ w \end{bmatrix} := \{x \in \mathbb{R}^n : x = V\mu \wedge \mu \geq 0 \wedge w^T \mu = 1\}$$

2. Illustration

Here $n = 2$, $k = 3$, and $\ell = 4$.



3. Tasks

- RR. Removing redundancy: if j is the number of non-redundant constraints (or vectors), this requires solving k (or ℓ) linear programming problems of size $n \times j$
- EN. Moving between H- and V-representations: done using vertex/facet enumeration algorithms; polynomial in n , k , and ℓ .
- PJ. Projection on a lower-dimensional space: easy with V-representations, hard with H-representations.
- IS. Intersection: easy with H-representations, hard with V-representations.



1. Formalization

Any multi-objective linear program (MOLP) can be put in the following form:

$$\begin{aligned} & \text{maximize } y = Cx, \\ & \text{subject to } Ax \leq b \text{ and } x \geq 0 \end{aligned}$$

3. Tasks

- Main computational tasks in non-decreasing order of complexity:
- M1. Finding \hat{y} .
 - M2. Finding \check{y} .
 - M3. Finding $\text{ext } \mathcal{Y}^*$ and characterizing \mathcal{Y}^* .
 - M4. Finding $\text{ext } \mathcal{X}^*$.
 - M5. Characterizing \mathcal{X}^* .

2. Illustration

Here $m = n = 2$ and $k = 4$.

