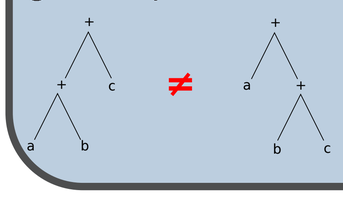
Numerical Accuracy Improvement of Programs: Principles and Experiments

Nasrine Damouche [1], Matthieu Martel [1] & Alexandre Chapoutot [2]

[1] **University of Perpignan** (Laboratoire LAMPS), 52 av Paul Alduy, 66860 PERPIGNAN Cedex 9
[2] **Paris-Saclay University** (U2IS, ENSTA ParisTech,), 828 bd des Maréchaux, 91762 Palaiseau cedex France
nasrine.damouche@univ-perp.fr, matthieu.martel@univ-perp.fr, alexandre.chapoutot@ensta-paristech.fr

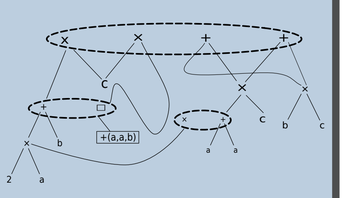
Problem

In general, the correctness of numerical computations of programs [4] based on floating-point arithmetic is not intuitive [1] and developers hope to compute an accurate result without guaranty. To solve this problem, we proceed by automatic source to source transformation of programs to improve their numerical accuracy.



Transforming Expressions

For automatic transformation of single arithmetic expressions, several techniques have already been proposed [5, 6]. Among them, the APEG (Abstract Program Expression Graph) introduce a new intermediary representation that manipulates in a single data structure a large set of equivalent arithmetic expressions [5]. Basically, APEG do not duplicate the common parts of the syntactic trees of mathematically equivalent expressions. The APEG corresponding to the expression $e = (a + a) + b * c$ is displayed hereafter. An expression is built from an APEG by selecting one operator in each dotted box and one arbitrary parsing for each box (e.g. $\boxed{+(a,a,b)}$).



Transforming Commands

The transformation uses an environment δ which maps identifiers to formal expressions, a black list β of identifiers and a target variable μ to be optimized. Programs are assumed to be given in SSA form and Φ-nodes are associated to conditionals and loops.

Assignments $c \equiv id = e$

- If the following conditions are satisfied $\Rightarrow$ remove the assignment from the program and memorize $e[id \mapsto e]$ in δ.
  - the variables of e do not exist in δ,
  - $v \notin \beta$ and $v \neq \nu$,
- Otherwise, we inline the variables saved in δ in the concerned expression e and we transform the resulting expression.

Sequences $c \equiv c_1; c_2$

- If c_1 or c_2 is nop $\Rightarrow$ rewrite the other member of the sequence,
- Otherwise, rewrite both members of the sequence.

Conditionals $c \equiv \mathsf{if}_\Phi\ e\ \mathsf{then}\ c_1\ \mathsf{else}\ c_2$

- If e is true $\Rightarrow$ rewrite c_1,
- If e is false $\Rightarrow$ rewrite c_2,
- If e is statically unknown $\Rightarrow$ rewrite both branches of the conditionals,
- If the variables of e have been removed from the program $\Rightarrow$ re-insert them into the source code and then rewrite it (the variables of e are inserted in the black list β).

While Loops $c \equiv \mathsf{while}_\Phi\ e\ \mathsf{do}\ c$

- Rewrite the body of the loop,
- re-inject the variables discarded into the program and then rewrite it (again the concerned variables are added to β).

Experiments and Results

To illustrate how we transform programs, we have taken an example from robotics. It computes the position of a two wheeled robot by odometry [3]. Note that s_l and s_r are the instantaneous rotation speeds of the left and right wheels, c is the circumference of the wheels and l the length of its axle. The left Figure represents the initial program of the odometry while the right one gives the transformed code.

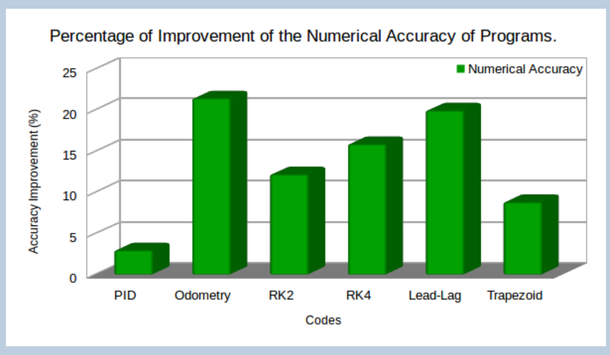
```
sl = [0.52,0.53];
sr = 0.785398163397; c = 12.34;
theta = t = x = y = 0.0 ; inv_l = 0.1 ;
while (t < 100.0 ) do {
  delta_dl = (c * sl);
  delta_dr = (c * sr);
  delta_d = ((delta_dl + delta_dr) * 0.5);
  delta_theta = ((delta_dr - delta_dl) * inv_l);
  arg = (theta + (delta_theta * 0.5));
  cos = (1.0 - ((arg * arg) * 0.5))
      + ((((arg * arg)* arg)* arg) / 24.0);
  x = (x + (delta_d * cos));
  sin = (arg - (((arg * arg)* arg)/6.0))
      + (((((arg * arg)* arg)* arg)* arg)/120.0);
  y = (y + (delta_d * sin));
  theta = (theta + delta_theta);
  t = (t + 0.1);}
```

Listing of the initial Odometry program.

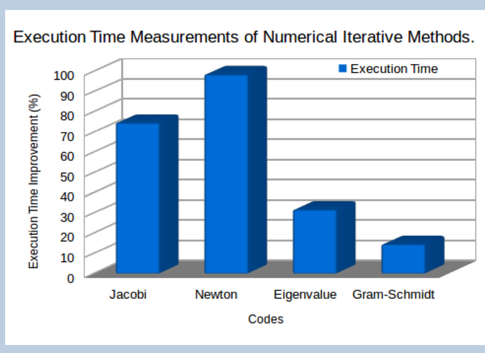
```
sl = [0.52,0.53]; theta = t = x = y = 0.0 ;
while (t < 100.0) do {
  TMP_6 = (0.1 * (0.5 * (9.691813336318980 - (12.34 * sl))));
  TMP_23 = ((theta + (((9.691813336318980 - (sl * 12.34))
         * 0.1) * 0.5)) * (theta + (((  9.691813336318980
         - (sl * 12.34)) * 0.1) * 0.5)));
  TMP_25 = ((theta + TMP_6) * (theta + TMP_6)) * (theta
         + (((9.691813336318980 - (sl * 12.34)) * 0.1) * 0.5));
  TMP_26 = (theta + TMP_6) ;
  x = ((0.5 * (((1.0 - (TMP_23 * 0.5)) + ((TMP_25 * TMP_26)
    / 24.0)) * ((12.34 * sl) + 9.691813336318980))) + x);
  TMP_27 = ((TMP_26 * TMP_26) * (theta
         + (((9.691813336318980 - (sl * 12.34)) * 0.1) * 0.5)));
  TMP_29 = (((TMP_26 * TMP_26) * TMP_26) * (theta
         + (((9.691813336318980 - (sl * 12.34)) * 0.1) * 0.5)));
  y = (((9.691813336318980 + (12.34 * sl)) * (((TMP_26
    - (TMP_27 / 6.0)) + ((TMP_29 * TMP_26) / 120.0)) * 0.5)) + y);
  theta = (theta + (0.1 * (9.691813336318980 - (12.34 * sl))));
  t = t + 0.1;}
```

Listing of the transformed Odometry program.

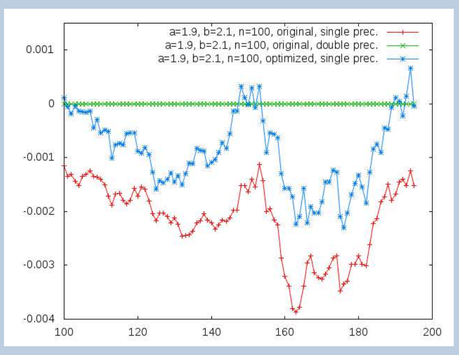
For the demonstration of its efficiency, we evaluate through our tool a set of programs in the domains of robotics, avionics, mathematics, etc. The experimental results show the efficiency of our approach to optimize the numerical accuracy. We compare the initial and the new error of each of programs (see Figure (a)). We have also demonstrated [2] that improving the numerical accuracy accelerates the convergence time of numerical iterative algorithms (see Figure (b)). Finally, we have compared initial codes working in double precision with optimized codes in simple precision (see figure (c)).



(a)



(b)



(c)

Conclusion and Perspectives

- Improve numerical accuracy of programs by up to 20%,
- Accelerate convergence of numerical iterative methods up to 15%,
- Transformed programs in single precision are close to source programs with double precision.
- Extend this work to deal with other or more complicated structures of programs such that functions and arrays.
- Validate that programs transformed by our tool are equivalent to the original programs using the Coq proof assistant.
- Optimize the numerical accuracy of parallel programs.

References

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