# An optimality property of an approximated solution computed by the Hessenberg method

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**Abstract.** We revisit the implementation of the Krylov subspace method based on the Hessenberg process for general linear operator equations. It is established that at each step, the computed approximate solution can be regarded by the corresponding approach as the minimizer of a certain norm of residual corresponding to the obtained approximate solution of the system. Test problems are numerically examined for solving tensor equations with a cosine transform product arising from image restoration to compare the performance of the Krylov subspace methods in conjunction with the Tikhonov regularization technique based on Hessenberg and Arnoldi processes.

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#### 1. Introduction

A multidimensional array of data is called a tensor whose modes stand for its number of indices. Throughout this paper, vectors and matrices are denoted by lowercase and capital letters, respectively, and tensors are represented by Euler script.

For the sake of generality, we consider the following linear operator equation:

$$\mathcal{F}(\mathfrak{X}) = \mathfrak{G},\tag{1}$$

where  $\mathcal{F}(\cdot)$  is a given linear operator from  $\mathbb{R}^{I_1 \times I_2 \times \cdots \times I_N}$  onto  $\mathbb{R}^{I_1 \times I_2 \times \cdots \times I_N}$ . The tensor equation in the form (1) incorporates several classes of tensor equations recently mentioned in the literature, including multilinear systems [6, 9, 19], the Sylvester matrix equation [4, 5, 7], the Stein tensor equation [4], etc. Basically, special cases of equation (1) appear in numerous areas such as a Markov process [8], physics [9], and numerical discretization of (high order) partial differential equations [4, 5, 7, 19] from engineering problems.

Based on the Arnoldi process, several variants of Krylov subspace methods have been developed in the literature for solving systems in the form (1), see [4, 5, 9, 10] and references therein. To be more precise, let us consider the Sylvester tensor equation (STE) as a special case of (1) and review some of recently proposed methods to solve STEs. To this end, first, we need to present the definitions of the mode-nproduct [22].

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**Definition 1.** The n-mode (matrix) product of a tensor  $\mathfrak{X} \in \mathbb{R}^{I_1 \times I_2 \times \cdots \times I_N}$  with a matrix  $U \in \mathbb{R}^{J \times I_n}$  is denoted by  $\mathfrak{X} \times_n U$  and is of size

$$I_1 \times \cdots \times I_{n-1} \times J \times I_{n+1} \times \cdots \times I_N,$$

and its elements are defined as follows:

$$(\mathfrak{X} \times_n U)_{i_1 \cdots i_{n-1} j i_{n+1} \cdots i_N} = \sum_{i_n=1}^{I_n} x_{i_1 i_2 \cdots i_N} u_{j i_n}.$$

Consider the STE as follows:

$$\mathfrak{X} \times_1 A^{(1)} + \mathfrak{X} \times_2 A^{(2)} + \dots + \mathfrak{X} \times_N A^{(N)} = \mathcal{D}, \qquad (2)$$

where the right-hand side tensor  $\mathcal{D} \in \mathbb{R}^{I_1 \times I_2 \times \cdots \times I_N}$  and coefficient matrices  $A^{(n)} \in \mathbb{R}^{I_n \times I_n}$   $(n = 1, 2, \ldots, N)$  are known, and  $\mathcal{X} \in \mathbb{R}^{I_1 \times I_2 \times \cdots \times I_N}$  is unknown. In the literature, several variants of the Krylov subspace methods were proposed for solving the above STE, see [5, 7, 10, 23, 26] and the references therein. In particular, Kressner and Tobler [23] applied Krylov subspace methods based on the (extended) Arnoldi process for the case that the right-hand side  $\mathcal{D}$  is a tensor of low rank. For the case when  $\mathcal{D}$  does not necessarily have a low rank, Chen and Lu [10] developed the generalized minimal residual (GMRES) method in a tensor framework. The tensor form of the full orthogonalization method (FOM) was presented in [6]. STEs with dense coefficient matrices  $A^{(1)}, A^{(2)}, \ldots, A^{(N)}$  can possibly arise from discretization of three-dimensional partial differential equations by spectral methods [24, 25]. In [26], it is observed that using the Hessenberg process instead of the Arnoldi process can lead to a computationally cheaper Krylov subspace method when the coefficient matrices in the STE are dense.

It is known that replacing the Hessenberg process by the Arnoldi process can lead to cost-effective Krylov subspace methods, see [18, 29]. This fact motivated several researchers to extend Krylov subspace methods based on the Hessenberg process for solving different types of linear operator equations in the form (1). For instance, the block Changing Minimal Residual method based on the Hessenberg (CMRH) method was proposed in [1] to solve linear systems of the form AX = B, where A is nonsymmetric. The weighted and flexible versions of block CMRH were also presented in [2]. Gu et al. [13] proposed a restarted Hessenberg method to solve shifted nonsymmetric linear systems. In [14], the restarted CMRH process was presented for solving multi-shifted linear systems with non-Hermitian coefficient matrices. Recently, a Hessenberg-type method was applied for the solution of PageRank problems, see [15]. Brief discussions are included in Appendix A to recall Hessenberg and Arnoldi processes associated with linear operator  $\mathcal{F}(\cdot)$  in tensor equation (1) and compare their computational costs.

The remainder of this paper is organized as follows: In Section 2, we briefly explain how the Hessenberg method is used for solving (1). The optimality property of an approximate solution obtained by the Hessenberg method is established in Section 3. Numerical experiments are reported in Section 4 to disclose comparison results between Hessenberg and Arnoldi processes in conjunction with the Tikhonov regularization technique for a class of tensor equations arising from image restoration. We finish the paper by some concluding remarks in Section 5.

### 2. An overview of the Hessenberg method

In this section, we briefly review the implementation of the Hessenberg method for (1). To do so, we need to recall the inner product between two tensors and its induced norm. The inner product between two same size tensors  $\mathcal{X}$  and  $\mathcal{Y}$  in  $\mathbb{R}^{I_1 \times I_2 \times \cdots \times I_N}$  is given by

$$\langle \mathfrak{X}, \mathfrak{Y} \rangle = \sum_{i_1=1}^{I_1} \sum_{i_2=1}^{I_2} \dots \sum_{i_N=1}^{I_N} x_{i_1 i_2 \dots i_N} y_{i_1 i_2 \dots i_N}.$$
 (3)

Given a tensor  $\mathfrak{X} \in \mathbb{R}^{I_1 \times I_2 \times \cdots \times I_N}$ , the induced norm from the above inner product is defined by

$$\|\mathfrak{X}\|^2 = \sum_{i_1=1}^{I_1} \sum_{i_2=1}^{I_2} \cdots \sum_{i_N=1}^{I_N} x_{i_1 i_2 \cdots i_N}^2.$$

Corresponding to the tensor  $\mathfrak{X} \in \mathbb{R}^{I_1 \times I_2 \times \cdots \times I_N}$ , frontal slices or column tensors of  $\mathfrak{X}$  have the following form:

$$\underbrace{\mathfrak{X}_{::\cdots:k}}_{(N-1)-\text{times}} \in \mathbb{R}^{I_1 \times I_2 \times \cdots \times I_{N-1}}, \quad k = 1, 2, \dots, I_N;$$

for further details, see [22]. When  $\mathfrak{X}$  is a tensor of order three, we also use the notation  $\mathfrak{X}^{(k)}$  to denote its k-th frontal slice.

Constructing iterative schemes based on the Hessenberg process for solving (1) follows from a similar strategy used in [26] and related details are omitted here. Basically, one can develop the Hessenberg method by constructing the basis  $\{\mathcal{V}_1, \mathcal{V}_2, \ldots, \mathcal{V}_m\}$  for the following Krylov subspace:

$$\mathcal{K}_m(\mathcal{F}, \mathcal{R}_0) = \operatorname{span}\{\mathcal{R}_0, \mathcal{F}(\mathcal{R}_0), \dots, \mathcal{F}^{m-1}(\mathcal{R}_0)\},\tag{4}$$

using Algorithm 1 in Appendix A such that

$$\langle \mathcal{V}_{i+1}, \mathcal{Y}_j \rangle = 0, \quad \text{for} \quad j = 1, 2, \dots, i,$$

in which the linearly independent tensors  $\mathcal{Y}_i$ s of order  $I_1 \times I_2 \times \cdots \times I_N$  are available,  $\mathcal{R}_0 = \mathcal{G} - \mathcal{F}(\mathcal{X}_0)$  and the initial guess  $\mathcal{X}_0 \in \mathbb{R}^{I_1 \times I_2 \times \cdots \times I_N}$  is given.

In what follows, we define

$$\bar{H}_m := \begin{pmatrix} H_m \\ e_m^\top h_{m+1,m} \end{pmatrix},$$

where the (i, j)-th entry of  $H_m$  is denoted by  $h_{ij}$  computed in lines 5 and 8 of Algorithm 1.

Suppose that  $\tilde{\mathcal{V}}_m$  and  $\tilde{\mathcal{Y}}_m$  are (N+1)-mode tensors with column tensors  $\mathcal{V}_i$ s and  $\mathcal{Y}_i$ s for  $i = 1, 2, \ldots, m$ . The following theorem is useful to derive Krylov subspace methods based on the Hessenberg process for linear operator equations in the form (1). The proof of theorem follows from similar strategies used in [5, 26].

**Theorem 1.** Let  $\tilde{W}_m$  be the (N+1)-mode tensor with column tensors  $W_j := \mathcal{F}(V_j)$  for j = 1, ..., m. Then the following statements hold:

$$\tilde{\mathcal{W}}_m = \tilde{\mathcal{V}}_{m+1} \times_{(N+1)} \bar{H}_m^\top,\tag{5}$$

$$\tilde{\mathcal{W}}_m = \tilde{\mathcal{V}}_m \times_{(N+1)} H_m^\top + h_{m+1,m} \mathcal{Z} \times_{(N+1)} E_m, \tag{6}$$

in which  $\mathcal{Z}$  is an (N+1)-mode tensor with "m" column tensors  $0, \ldots, 0, \mathcal{V}_{m+1}$  and  $E_m$  is an  $m \times m$  matrix of the form  $E_m = [0, \ldots, 0, e_m]$ , where  $e_m$  is the m-th column of the identity matrix of order m.

Let  $\mathcal{V}_1, \mathcal{V}_2, \ldots, \mathcal{V}_m$  be a basis for  $\mathcal{K}_m(\mathcal{F}, \mathcal{R}_0)$  produced via Algorithm 1. The *m*-th approximate solution  $\mathcal{X}_m$  is determined such that

$$\mathfrak{X}_m \in \mathfrak{X}_0 + \mathcal{K}_m(\mathcal{F}, \mathfrak{R}_0),$$

which implies that

$$\mathfrak{X}_m = \mathfrak{X}_0 + \sum_{i=1}^m \mathfrak{V}_i y_m^{(i)}.$$
(7)

The following orthogonality conditions in the Hessenberg method are imposed:

$$\langle \mathfrak{R}_m, \mathfrak{Y}_i \rangle = 0, \quad \text{for} \quad i = 1, 2, \dots, m,$$
(8)

to obtain the unknown vector  $y_m = (y_m^{(1)}; y_m^{(2)}; \ldots; y_m^{(m)})$  where  $\mathcal{R}_m = \mathcal{G} - \mathcal{F}(\mathfrak{X}_m)$ and the MATLAB notation  $(w_1; w_2; \ldots; w_m)$  represents the vector  $(w_1, w_2, \ldots, w_m)^\top$ . Using Theorem 1 it is not difficult to verify that  $y_m$  satisfies

$$H_m y_m = \beta e_1,\tag{9}$$

with  $\beta = \langle \mathcal{R}_0, \mathcal{Y}_1 \rangle$ . It can be verified that

$$\mathfrak{R}_m = \mathfrak{G} - \mathcal{F}(\mathfrak{X}_m) = -h_{m+1,m} y_m^{(m)} \mathcal{V}_{m+1},$$

for more details, see [26].

## 3. An optimality property of the Hessenberg method

Let  $\mathcal{F} : \mathbb{R}^{I_1 \times I_2 \times \cdots \times I_N} \to \mathbb{R}^{I_1 \times I_2 \times \cdots \times I_N}$  be a given arbitrary invertible linear operator, i.e.,  $\mathcal{F}(\mathcal{X}) = 0$  implies  $\mathcal{X} = 0$ . In this section, we show that the computed approximate solution by the Hessenberg method at each step satisfies an optimality property. To this end, first we need to recall a special case of the contracted tensor product.

**Definition 2** (see [5]). The  $\boxtimes^N$  product between N-mode tensors  $\mathfrak{X} \in \mathbb{R}^{I_1 \times \cdots \times I_{N-1} \times I_N}$ and  $\mathfrak{Y} \in \mathbb{R}^{I_1 \times \cdots \times I_{N-1} \times \tilde{I}_N}$  is defined as an  $I_N \times \tilde{I}_N$  matrix whose (i, j)-th entry is

$$[\mathfrak{X} \boxtimes^N \mathfrak{Y}]_{ij} = tr(\mathfrak{X}_{::\dots:i} \boxtimes^{N-1} \mathfrak{Y}_{::\dots:j}), \qquad N = 3, 4, \dots,$$

where

$$\mathfrak{X} \boxtimes^2 \mathfrak{Y} = \mathfrak{X}^\top \mathfrak{Y}, \qquad \mathfrak{X} \in \mathbb{R}^{I_1 \times I_2}, \mathfrak{Y} \in \mathbb{R}^{I_1 \times I_2}$$

We comment that the  $\boxtimes^N$  product between  $\mathfrak{X}$  and  $\mathfrak{Y}$  is a reformulation of a special case of the contracted product [11]. As pointed out in [5], it can be verified that

$$\langle \mathfrak{X}, \mathfrak{Y} \rangle = \operatorname{tr}(\mathfrak{X} \boxtimes^N \mathfrak{Y}), \qquad N = 2, 3, \dots,$$

$$(10)$$

for  $\mathfrak{X}, \mathfrak{Y} \in \mathbb{R}^{I_1 \times I_2 \times \cdots \times I_N}$ .

Now we define a new inner product and its associated norm as well in order to establish the following proposition showing that the computed approximate solution by the Hessenberg method is the minimizer of a residual corresponding to the approximate solution of (1).

**Definition 3.** Let  $\tilde{\mathbb{Y}}_k$  be an (N + 1)-mode tensor with frontal slices  $\mathbb{Y}_i$  for  $i = 1, 2, \ldots, m$  and let  $\epsilon > 0$  be given. For  $\mathfrak{X}, \mathfrak{Z} \in \mathbb{R}^{I_1 \times I_2 \times \cdots \times I_N}$ , we define the following inner product:

$$\langle \mathfrak{X}, \mathfrak{Z} \rangle_{\tilde{\mathfrak{Y}}_{k}, \epsilon} = \left\langle \tilde{\mathfrak{Y}}_{k} \boxtimes^{(N+1)} \mathfrak{X}, \tilde{\mathfrak{Y}}_{k} \boxtimes^{(N+1)} \mathfrak{Z} \right\rangle + \epsilon \left\langle \mathfrak{X}, \mathfrak{Z} \right\rangle.$$
(11)

The corresponding tensor norm is given by  $\|X\|^2_{\tilde{\mathfrak{Y}}_{k},\epsilon} = \langle \mathfrak{X}, \mathfrak{X} \rangle_{\tilde{\mathfrak{Y}}_{k},\epsilon}.$ 

We add the following remark to the previous definition to clarify the fact that the bilinear form (11) is an inner product.

**Remark 1.** Considering equality (10), one can see that

$$\left<\tilde{\mathcal{Y}}_k\boxtimes^{(N+1)}\mathfrak{X}, \tilde{\mathcal{Y}}_k\boxtimes^{(N+1)}\mathfrak{Z}\right> = \left< w_x, w_z \right>$$

where  $w_x = (\langle \mathfrak{Y}_1, \mathfrak{X} \rangle; \langle \mathfrak{Y}_2, \mathfrak{X} \rangle; \ldots; \langle \mathfrak{Y}_m, \mathfrak{X} \rangle)$  and  $w_z = (\langle \mathfrak{Y}_1, \mathfrak{Z} \rangle; \langle \mathfrak{Y}_2, \mathfrak{Z} \rangle; \ldots; \langle \mathfrak{Y}_m, \mathfrak{Z} \rangle)$ . Therefore, the bilinear form (11) is basically the summation of two inner products. Since  $\epsilon > 0$ , it is not difficult to verify that the bilinear form (11) is indeed an inner product.

**Proposition 1.** Assume that the linear operator  $\mathcal{F}(\cdot)$  is invertible, i.e.,  $\mathcal{F}(\mathfrak{X}) = 0$  implies  $\mathfrak{X} = 0$ . Let  $\mathfrak{X}_k$  be the k-th approximate solution of  $\mathcal{F}(\mathfrak{X}) = \mathfrak{G}$  obtained after implementing the Hessenberg method. Then, there exists  $\hat{\epsilon} > 0$  such that

$$\left\| \mathfrak{G} - \mathcal{F}(\mathfrak{X}_k) \right\|_{\tilde{\mathfrak{Y}}_{k,\epsilon}} < \left\| \mathfrak{G} - \mathcal{F}(\hat{\mathfrak{X}}) \right\|_{\tilde{\mathfrak{Y}}_{k,\epsilon}} \quad for \quad 0 < \epsilon < \hat{\epsilon}, \tag{12}$$

for any  $\hat{\mathfrak{X}} \in \mathfrak{X}_0 + \mathcal{K}_k(\mathcal{F}, \mathfrak{R}_0)$  provided that  $\tilde{\mathfrak{Y}}_k \boxtimes^{(N+1)} \mathcal{F}(\mathfrak{X}_k - \hat{\mathfrak{X}}) \neq 0$ . Here the frontal slices of the (N+1)-mode tensor  $\tilde{\mathfrak{Y}}_k$  are given by  $\mathfrak{Y}_1, \mathfrak{Y}_2, \ldots, \mathfrak{Y}_k$  such that

$$\langle \mathfrak{R}_k, \mathfrak{Y}_i \rangle = 0, \quad for \quad i = 1, 2, \dots, k$$

in which  $\Re_k = \Im - \mathcal{F}(\mathfrak{X}_k)$ .

**Proof.** Let  $\hat{\mathfrak{X}} \in \mathfrak{X}_0 + \mathcal{K}_k(\mathcal{F}, \mathfrak{R}_0)$  and  $\tilde{\mathfrak{Y}}_k \boxtimes^{(N+1)} \mathcal{F}(\mathfrak{X}_k - \hat{\mathfrak{X}}) \neq 0$ . It is not difficult to verify that

$$\begin{split} \left\| \mathfrak{G} - \mathcal{F}(\hat{\mathfrak{X}}) \right\|_{\tilde{\mathfrak{Y}}_{k,\epsilon}}^{2} &= \left\langle \mathfrak{G} - \mathcal{F}(\hat{\mathfrak{X}}), \mathfrak{G} - \mathcal{F}(\hat{\mathfrak{X}}) \right\rangle_{\tilde{\mathfrak{Y}}_{k,\epsilon}} \\ &= \left\langle \mathfrak{R}_{k} + \mathcal{F}(\mathfrak{X}_{k} - \hat{\mathfrak{X}}), \mathfrak{R}_{k} + \mathcal{F}(\mathfrak{X}_{k} - \hat{\mathfrak{X}}) \right\rangle_{\tilde{\mathfrak{Y}}_{k,\epsilon}} \\ &= \left\| \mathfrak{R}_{k} \right\|_{\tilde{\mathfrak{Y}}_{k,\epsilon}}^{2} + 2 \left\langle \mathfrak{R}_{k}, \mathcal{F}(\mathfrak{X}_{k} - \hat{\mathfrak{X}}) \right\rangle_{\tilde{\mathfrak{Y}}_{k,\epsilon}} + \left\| \mathcal{F}(\mathfrak{X}_{k} - \hat{\mathfrak{X}}) \right\|_{\tilde{\mathfrak{Y}}_{k,\epsilon}}^{2}. \quad (13) \end{split}$$

Notice that  $\mathcal{F}(\hat{\mathcal{X}} - \mathcal{X}_k) \neq 0$ , hence, the invertibility of  $\mathcal{F}(\cdot)$  implies that  $\mathcal{F}(\hat{\mathcal{X}}) \neq \mathcal{F}(\mathcal{X}_k)$ . From (13), one can see that

$$\begin{split} \left\| \mathfrak{G} - \mathcal{F}(\hat{\mathfrak{X}}) \right\|_{\tilde{\mathfrak{Y}}_{k,\epsilon}}^{2} &= \left\| \mathfrak{R}_{k} \right\|_{\tilde{\mathfrak{Y}}_{k,\epsilon}}^{2} + 2 \left( \left\langle \tilde{\mathfrak{Y}}_{k} \boxtimes^{(N+1)} \mathfrak{R}_{k}, \tilde{\mathfrak{Y}}_{k} \boxtimes^{(N+1)} \mathcal{F}(\mathfrak{X}_{k} - \hat{\mathfrak{X}}) \right\rangle \\ &+ \epsilon \left\langle \mathfrak{R}_{k}, \mathcal{F}(\mathfrak{X}_{k} - \hat{\mathfrak{X}}) \right\rangle \right) + \left\| \mathcal{F}(\mathfrak{X}_{k} - \hat{\mathfrak{X}}) \right\|_{\tilde{\mathfrak{Y}}_{k,\epsilon}}^{2}. \end{split}$$

Invoking the fact that  $\tilde{\mathcal{Y}}_k \boxtimes^{(N+1)} \mathcal{R}_k$  is a zero vector, we obtain

$$\begin{aligned} \left\| \mathfrak{G} - \mathcal{F}(\hat{\mathfrak{X}}) \right\|_{\tilde{\mathfrak{Y}}_{k,\epsilon}}^{2} &= \left\| \mathfrak{R}_{k} \right\|_{\tilde{\mathfrak{Y}}_{k,\epsilon}}^{2} + \left\| \mathcal{F}(\mathfrak{X}_{k} - \hat{\mathfrak{X}}) \right\|_{\tilde{\mathfrak{Y}}_{k,\epsilon}}^{2} + 2\epsilon \left\langle \mathfrak{R}_{k}, \mathcal{F}(\mathfrak{X}_{k} - \hat{\mathfrak{X}}) \right\rangle \\ &> \left\| \mathfrak{R}_{k} \right\|_{\tilde{\mathfrak{Y}}_{k,\epsilon}}^{2} + \left\| \mathcal{F}(\mathfrak{X}_{k} - \hat{\mathfrak{X}}) \right\|_{\tilde{\mathfrak{Y}}_{k,\epsilon}}^{2} - 2\epsilon \left| \left\langle \mathfrak{R}_{k}, \mathcal{F}(\mathfrak{X}_{k} - \hat{\mathfrak{X}}) \right\rangle \right|. \tag{14}$$

If  $\langle \mathcal{R}_k, \mathcal{F}(\mathcal{X}_k - \hat{\mathcal{X}}) \rangle = 0$ , the assertion follows from the above inequality for any  $\epsilon > 0$ . Consequently, without loss of generality, we assume that  $\langle \mathcal{R}_k, \mathcal{F}(\mathcal{X}_k - \hat{\mathcal{X}}) \rangle \neq 0$  and define

$$\hat{\epsilon} := \frac{\left\| \tilde{\mathfrak{Y}}_k \boxtimes^{(N+1)} \mathcal{F}(\mathfrak{X}_k - \hat{\mathfrak{X}}) \right\|^2}{2 \left| \left\langle \mathfrak{R}_k, \mathcal{F}(\mathfrak{X}_k - \hat{\mathfrak{X}}) \right\rangle \right|}.$$
(15)

In view of (14), for any  $\epsilon < \hat{\epsilon}$ , we have

$$\begin{split} \left\| \mathfrak{G} - \mathcal{F}(\hat{\mathfrak{X}}) \right\|_{\tilde{\mathfrak{Y}}_{k,\epsilon}}^{2} &> \left\| \mathfrak{R}_{k} \right\|_{\tilde{\mathfrak{Y}}_{k,\epsilon}}^{2} + \left\| \mathcal{F}(\mathfrak{X}_{k} - \hat{\mathfrak{X}}) \right\|_{\tilde{\mathfrak{Y}}_{k,\epsilon}}^{2} - 2\hat{\epsilon} \left| \left\langle \mathfrak{R}_{k}, \mathcal{F}(\mathfrak{X}_{k} - \hat{\mathfrak{X}}) \right\rangle \right| \\ &= \left\| \mathfrak{R}_{k} \right\|_{\tilde{\mathfrak{Y}}_{k,\epsilon}}^{2} + \epsilon \left\| \mathcal{F}(\mathfrak{X}_{k} - \hat{\mathfrak{X}}) \right\|^{2} \\ &> \left\| \mathfrak{R}_{k} \right\|_{\tilde{\mathfrak{Y}}_{k,\epsilon}}^{2} = \left\| \mathfrak{G} - \mathcal{F}(\mathfrak{X}_{k}) \right\|_{\tilde{\mathfrak{Y}}_{k,\epsilon}}^{2}, \end{split}$$

which completes the proof.

We end this part with the following remark on choosing  $\epsilon$ .

**Remark 2.** Assume that the assumptions of Proposition 1 hold. In view of (15) and the Cauchy–Schwarz inequality, one may set

$$\hat{\epsilon} = \frac{1}{2} \cdot \min_{0 \neq \mathcal{Z} \in \mathcal{K}_k(\mathcal{F}, \mathcal{R}_0)} \left\{ \frac{\|\mathcal{W}(\mathcal{Z})\|}{\|\mathcal{R}_k\|} \mid \mathcal{R}_k \neq 0 \quad and \quad \mathcal{W}(\mathcal{Z}) \neq 0 \right\}$$

to eliminate the dependency of  $\hat{\epsilon}$  on  $\hat{\mathfrak{X}}$  in the statement of Proposition 1 where  $\mathcal{W}(\mathfrak{Z}) := \tilde{\mathfrak{Y}}_k \boxtimes^{(N+1)} \mathcal{F}(\mathfrak{Z})$ . In fact, the k-th approximate solution obtained via the Hessenberg method satisfies the optimality property (12) for any  $0 < \epsilon < \hat{\epsilon}$ .

# 4. An application from image processing

Developing efficient image deblurring methods is an active area of research. For iterative methods based on the Krylov subspace, one can refer to [4, 7, 12, 26, 30,

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31, 32] and the references therein. For example, Guo et al. [16] recently developed a three-dimensional fractional total variation– based model for a three-dimensional image deblurring problem.

In this section, we aim to experimentally illustrate the performance of the Hessenberg method in the context of image deblurring. For this application, the implementation of the Hessenberg method has been already considered in [26], where equation (1) is reduced to the Sylvester tensor equation. Here, we implement the method for solving an alternative tensor equation. To this end, before presenting numerical results, we need to review some preliminaries.

#### 4.1. Basic concepts

To present our mentioned tensor equation in reported numerical experiments, we need to recall the definition of the  $*_c$ -product from [20].

**Definition 4** (\**<sub>c</sub>*-product). Let  $\mathcal{A} \in \mathbb{R}^{m \times \ell \times n}$  and  $\mathcal{B} \in \mathbb{R}^{\ell \times p \times n}$ . The tensor-tensor product  $\mathcal{C} = \mathcal{A} *_{c} \mathcal{B}$  is of size  $m \times p \times n$  such that

$$\hat{\mathcal{C}}^{(i)} := \hat{\mathcal{A}}^{(i)} \hat{\mathcal{B}}^{(i)}, \qquad \text{for } i = 1, \dots, n,$$

where  $\hat{\mathcal{A}} = \mathcal{A} \times_3 M$ ,  $\hat{\mathcal{B}} = \mathcal{B} \times_3 M$  and  $\mathcal{C} = \hat{\mathcal{C}} \times_3 M^{-1}$ . The matrix  $M = W^{-1}C(I+Z)$ can be computed in MATLAB using

$$C = dct(eye(n)), W = diag(C(:,1)), Z = diag(ones(n-1,1),1)$$

We consider the following tensor equation:

$$\mathcal{A} *_c \mathfrak{X} = \mathfrak{G},\tag{16}$$

where tensor  $\mathcal{A} \in \mathbb{R}^{\ell \times \ell \times m}$  and the right-hand side  $\mathcal{G} \in \mathbb{R}^{\ell \times p \times m}$  are given and  $\mathfrak{X} \in \mathbb{R}^{\ell \times p \times m}$  is the unknown tensor. The tensor problem (16) may appear in engineering, signal processing, and image and video data processing problems, see [12, 20, 21, 27, 28].

Notice that the linear operator  $\mathcal{F}(\cdot)$  takes the following form:

$$\begin{aligned} \mathcal{F} : \mathbb{R}^{\ell \times p \times m} &\to \mathbb{R}^{\ell \times p \times m} \\ \mathcal{X} &\mapsto \mathcal{F}(\mathcal{X}) := \mathcal{A} *_c \mathcal{X} \end{aligned}$$

In the sequel, we briefly explain the strategy for exploiting tensors to reformat a typical discrete model for image blurring; for further details, see [20]. Consider the linear system of equations

$$Bx = g, (17)$$

where *B* denotes the discrete blurring matrix of order  $n^2$  and *x* is the vectorized form of the image *X*. The right-hand side of the above equation contains an error *e* called "noise", i.e.,  $g = \hat{g} + e$ , where  $\hat{g}$  is the unknown noise-free unavailable right-hand side. Following the discussions in [20], one can reformulate equation (17) by

$$\mathcal{A} *_c \mathfrak{X} = \mathfrak{G}, \qquad \mathfrak{G} = \mathfrak{G} + \mathfrak{N}, \tag{18}$$

where the image X is obtained by "reshaping" the first column of  $\mathtt{mat}(\mathfrak{X})$ , and the right-hand side tensor  $\mathfrak{G}$  is obtained by  $\mathtt{ten}(g)$ , for the definition of  $\mathtt{mat}(\cdot)$  and  $\mathtt{ten}(\cdot)$ , see [20]. The tensor  $\mathfrak{G}$  in (18) is contaminated by a noise tensor  $\mathfrak{N}$  with normally distributed random entries with zero mean and scaled to correspond to a specific noise level  $\eta = \|\mathfrak{N}\| / \|\mathfrak{G}\|$ .

Let  $B = T \otimes T$ , T is a Toeplitz matrix representing a Gaussian blur. We created  $T \in \mathbb{R}^{n \times n}$  in MATLAB as follows:

$$T = \frac{1}{\sqrt{2\pi\sigma^2}} \cdot \operatorname{toeplitz}\left(\mathbf{z}\right),\tag{19}$$

with

$$\mathtt{z} = \left[ \mathtt{exp} \left( -([0:\mathtt{band}-1]^2_{\cdot})/(2\sigma^2) 
ight), \mathtt{zeros}(1,n-\mathtt{band}) 
ight],$$

where the band and  $\sigma$  are given for each of our test problems, and "toeplitz(·)" is a command in MATLAB. To generate the coefficient tensor  $\mathcal{A}$  in (18), we set  $\mathcal{A}^{(i)} = T(i,1)T$  for i = 1, 2, ..., n, recalling that  $\mathcal{A}^{(i)}$  denotes the *i*-th frontal slice of  $\mathcal{A}$ .

#### 4.2. Numerical experiments

All numerical experiments were computed using MATLAB version 9.9 (R2020b) running on an Intel Core i5 CPU at 2.50 GHz with 8 GB of memory using Tensor Toolbox [3].

The Tikhonov regularization technique consists of replacing the solution of (18) by the following minimization problem:

$$\min_{\mathfrak{X}\in\mathbb{R}^{\ell\times p\times m}}\left\{\|\mathcal{A}*_{c}\mathfrak{X}-\mathcal{G}\|^{2}+\mu\|\mathfrak{X}\|^{2}\right\},\$$

where  $\mu > 0$  is the regularization parameter. In what follows, we compare the performance of Hessenberg and Arnoldi methods in conjunction with the Tikhonov regularization method. The corresponding methods are called Hessenberg-Tikhonov and Arnoldi-Tikhonov methods, respectively. For more details on the implementation of the Hessenberg and Arnoldi methods in conjunction with the Tikhonov regularization method, we refer the readers to [6, 7, 26].

In Table 1, we report the total required number of iterations and consumed CPUtime (in seconds) under "Iter" and "CPU(s)", respectively. For more details, we also disclose the relative error

$$\operatorname{Err} := \frac{\|\mathfrak{X}_{\mu_k,k} - \hat{\mathfrak{X}}\|}{\|\hat{\mathfrak{X}}\|},$$

where  $\hat{X}$  denotes the exact solution of the problem with the error-free right-hand side tensor  $\hat{\mathcal{G}}$  associated with  $\mathcal{G}$ , and  $\mathcal{X}_{\mu_k,k}$  denotes the k-th computed approximation determined by the algorithms. We note that the deblurred images based on computed regularized solutions are obtained by reshaping the first column of  $\mathtt{mat}(\mathcal{X}_{\mu_k,k})$ . The regularization parameter  $\mu_k$  is determined by the discrepancy principle, for more details, see [17, Chapter 7]. The initial approximate solution in all experiments is the zero tensor and the iterations were terminated once a maximum 60 number of iterations is reached or the following condition holds:

$$\frac{\|\mathfrak{X}_{\mu_k,k} - \mathfrak{X}_{\mu_{k-1},k-1}\|}{\|\mathfrak{X}_{\mu_{k-1},k-1}\|} \le \tau,$$
(20)

for a user-specified value of the parameter  $\tau > 0$ .

In tables 1 and 2, the MATLAB function PSNR denotes the peak signal-to-noise ratio between the original and a blurred (or restored) image in decibels. The higher the PSNR, the better the quality of the deblurred image.

We examine the following two test problems to compare the performance of the Hessenberg-Tikhonov and Arnoldi-Tikhonov methods to restore an image contaminated by blur and noise.

**Example 1.** We use the blur operator obtained by (19) and set  $\tau = 5 \cdot 10^{-3}$  in (20). The results are reported for the following two cases:

**Case I.** band = 11 and  $\sigma = 4$ . The exact solution is the rice image from MATLAB. **Case II.** band = 16 and  $\sigma = 6$ . The exact solution is the airplane<sup>‡</sup> image.

Both gray-scale images are represented by an array of  $256 \times 256$  pixels. The original and blurred-noisy images are plotted in Figure 1. The blurred-noisy image G is obtained by reshaping the first column of mat(9). The obtained regularized solutions are shown in Figure 1 for the noise level  $\eta = 0.01$ .

Case I					
Noise level $(\eta)$	Method	Iter	$\mathrm{CPU}(\mathrm{s})$	Err	PSNR
0.01	Hessenberg-Tikhonov Arnoldi-Tikhonov	6 9	$0.6176 \\ 1.3402$	$\begin{array}{c} 1.3911 \cdot 10^{-1} \\ 1.8424 \cdot 10^{-1} \end{array}$	$\begin{array}{c} 23.4779 \\ 21.3066 \end{array}$
0.001	Hessenberg-Tikhonov Arnoldi-Tikhonov	8 12	$1.0822 \\ 2.1456$	$\begin{array}{c} 1.0984 \cdot 10^{-1} \\ 1.0965 \cdot 10^{-1} \end{array}$	$25.7960 \\ 25.8106$
Case II					
Noise level $(\eta)$	Method	Iter	$\mathrm{CPU}(\mathrm{s})$	Err	PSNR
0.01	Hessenberg-Tikhonov Arnoldi-Tikhonov	12 10	$2.1506 \\ 1.5554$	$\begin{array}{c} 8.1811 \cdot 10^{-2} \\ 8.5115 \cdot 10^{-2} \end{array}$	$\begin{array}{c} 24.0145 \\ 23.6695 \end{array}$
0.001	Hessenberg-Tikhonov Arnoldi-Tikhonov	6 10	$0.5611 \\ 1.4410$	$\begin{array}{c} 6.8164 \cdot 10^{-2} \\ 6.9903 \cdot 10^{-2} \end{array}$	$\begin{array}{c} 25.5971 \\ 25.3772 \end{array}$

Table 1: Comparison results for Example 1

We report the numerical results for Example 1 in Table 1. As observed, here both methods work well and determine suitable approximations for the exact solution for both noise levels. Overall, using the Hessenberg-Tikhonov method leads to

<sup>&</sup>lt;sup>†</sup>This image is available at: http://sipi.usc.edu/database/download.php?vol=misc&img=5.1.11



Figure 1: Original, noisy and restored images using Tikhonov regularization in conjunction with the Hessenberg and Arnoldi processes

better results than the Arnoldi-Tikhonov method in Case I. For the second case, the Hessenberg-Tikhonov method surpasses the Arnoldi-Tikhonov method for the level of noise  $\eta = 0.001$ . For the noise of level  $\eta = 0.01$ , the Hessenberg-Tikhonov method provides a slightly more accurate solution than the Arnoldi-Tikhonov method.

**Example 2.** The exact solution of this test example is the **boat**<sup>§</sup> image, which is represented by an array of  $512 \times 512$  pixels and displayed in Figure 2. We use the blur operator obtained by (19) with **band** = 3,  $\sigma = 4$ . The iterations are terminated as soon as the stopping criterion (20) is satisfied, where  $\tau = 4 \cdot 10^{-2}$ .

Results for this example are reported in Table 2. Both methods work well for both noise levels. The original, noisy and restored images are plotted in Figures 2 and 3. As seen, here, the Hessenberg-Tikhonov method consumes slightly less CPUtime (in seconds) in comparison with the Arnoldi-Tikhonov method. The provided approximate solutions by Hessenberg-Tikhonov are also a bit more accurate.

Noise level $(\eta)$	Method	Iter	CPU(s)	Err	PSNR
0.01	Hessenberg-Tikhonov Arnoldi-Tikhonov	$\frac{4}{4}$	$5.5581 \\ 5.6118$	$\begin{array}{c} 1.2472 \cdot 10^{-1} \\ 1.2915 \cdot 10^{-1} \end{array}$	$23.5330 \\ 23.2042$
0.001	Hessenberg-Tikhonov Arnoldi-Tikhonov	$4 \\ 4$	$5.5583 \\ 5.6160$	$\begin{array}{c} 1.0742 \cdot 10^{-1} \\ 1.1840 \cdot 10^{-1} \end{array}$	$24.8391 \\ 23.9633$

Table 2: Comparison results for Example 2

<sup>§</sup>This image is available at: https://sipi.usc.edu/database/download.php?vol=misc&img=boat.512

Optimality property of the Hessenberg method



Figure 2: Exact image (left) and contaminated image (right)



Figure 3: Restored images using Tikhonov regularization in conjunction with the Hessenberg and Arnoldi processes for noise of level 0.01

# 5. Conclusions

The Hessenberg method was considered to solve a general class of linear operator equations. It was shown that at each iterate the Hessenberg method produces an approximate solution satisfying an optimality property. The performance of the Hessenberg method had not been previously reported in the literature for solving the tensor equation  $\mathcal{A} *_c \mathfrak{X} = \mathfrak{G}$ . Therefore, two image restoration test problems in the above form were taken from [20]. Numerical comparison results were reported between the Hessenberg and Arnoldi processes in conjunction with the well-known

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Tikhonov regularization technique.

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## Appendix A.

In this part, we summarize the Hessenberg and Arnoldi processes. Brief discussions are also included to compare the computational costs of these processes. In addition, we recall the presented algorithm in [6] for finding the set of indices for which a tensor takes its maximum value in modulus.

Notice that the well-known Arnoldi process can be regarded as a special case of Algorithm 1 setting  $\mathcal{V}_i = \mathcal{Y}_i$  for  $i = 1, 2, \ldots, m$ . In practical implementation, in the Hessenberg process the tensor  $\mathcal{Y}_i$  is chosen as a tensor having only one nonzero entry equal to one. We comment that the iterative methods based on the Hessenberg process are applied with the pivoting strategy to avoid a possible breakdown, for further details, see [18, 29]. To this end, at each step of the Hessenberg process, Algorithm 2 is used for finding the set of indices corresponding to the maximum element (in modulus) of a tensor; for more details, see [26]. Basically, in this case, Algorithm 1 reduces to Algorithm 3.

We finish this part by comparing the number of required operations for the Arnoldi process and Algorithm 3 at each step of computing the new approximation for the solution of (1). Evidently, the differences between the number of operations in Arnoldi process and Algorithm 3 are in the requirement of using Algorithm 2, computing the values of  $\beta$  and  $h_{i,j}$  for  $i = 1, 2, \ldots, m + 1$  and  $j = 1, 2, \ldots, m$ . At each step, the total number of operations of the Arnoldi process is higher than Algorithm 3 due to more expensive computational costs of  $h_{i,j}$  in the Arnoldi process. Basically, the evaluation of each  $h_{i,j}$  corresponds to computing an inner product of the form (3). Consequently, lines 6 and 10 of Algorithm 3 demonstrate that each step of the Arnoldi process is more expensive than Algorithm 3. Notice that the cost of implementing Algorithm 2 is negligible, especially in the case when  $N \ll \max(I_1, I_2, \ldots, I_N)$ .

Algorithm 1 Hessenberg process. [26]

**Require:** Input tensor  $\mathcal{V}$  and scalar  $m \geq 1$  as the maximum allowed dimension of the Krylov subspace;

**Ensure:** The upper Hessenberg matrix  $\bar{H}_m = [h_{i,j}]_{(m+1)\times m}$  and (N+1)-mode tensor  $\tilde{\mathcal{V}}_m$  with the column tensors  $\mathcal{V}_1, \mathcal{V}_2, \ldots, \mathcal{V}_m$ .

1: Set  $\beta = \langle \mathcal{V}, \mathcal{Y}_1 \rangle$  and  $\mathcal{V}_1 = \mathcal{V}/\beta$ . 2: for j = 1, 2, ..., m do 3:  $\mathcal{W} = \mathcal{F}(\mathcal{V}_j)$ ; 4: for i = 1, 2, ..., j do 5:  $h_{i,j} = \langle \mathcal{Y}_i, \mathcal{W} \rangle$ ; 6:  $\mathcal{W} = \mathcal{W} - h_{i,j} \mathcal{V}_i$ ; 7: end for 8:  $h_{j+1,j} = \langle \mathcal{Y}_{j+1}, \mathcal{W} \rangle$ . If  $h_{j+1,j} = 0$ , then stop; 9:  $\mathcal{V}_{j+1} = \mathcal{W}/h_{j+1,j}$ ; 10: end for

Algorithm 2 Pivoting strategy for a  $\tau$ -mode tensor [6, Algorithm 1]

**Require:** Input a tensor  $\mathcal{R} \in \mathbb{R}^{I_1 \times I_2 \times \cdots \times I_{\tau}}$ ; **Ensure:** Index group  $(i_1, i_2, \dots, i_{\tau})$ ; 1:  $[\sim, j] = \max(|\operatorname{vec}(\mathcal{R})|)$ ; 2: **for**  $i = 1 : (\tau - 1)$  **do** 3:  $i_{\tau - i + 1} = \begin{bmatrix} \frac{j}{\prod_{\ell=1}^{\tau - i} I_\ell} \\ \frac{j}{\prod_{\ell=1}^{\tau - i} I_\ell} \end{bmatrix}$ ; 4:  $\ell = j - (i_{\tau - i + 1} - 1) \prod_{\ell=1}^{\tau - i} I_\ell$ ; 5:  $j = \ell$ ; 6: **end for** 7:  $i_1 = \ell - (i_2 - 1)I_1$ ;

Algorithm 3 Hessenberg\_BTF process with maximum strategy. [26]

**Require:** Input an  $I_1 \times I_2 \times \ldots \times I_N$ , tensor  $\mathcal{V}$  and the restart parameter m. **Ensure:** The upper Hessenberg matrix  $\bar{H}_m = [h_{i,j}]_{(m+1)\times m}$  and (N+1)-mode tensor  $\tilde{\mathcal{V}}_m$  with the column tensors  $\mathcal{V}_1, \mathcal{V}_2, \ldots, \mathcal{V}_m$ .

- 1: Determine triple  $(i_{1,0}, i_{2,0}, \ldots, i_{N,0})$  using Algorithm 2 for the input  $\mathcal{V}$ ;
- 2: Set  $\beta = \mathcal{V}_{i_{1,0},i_{2,0},\ldots,i_{N,0}}$ ;  $\mathcal{V}_1 = \mathcal{V}/\beta$ ; and  $p_{1,\eta} = i_{\eta,0}$  for  $\eta = 1, 2, \ldots, N$ ;
- 3: for j = 1, ..., m do
- 4:  $\mathcal{U} = \mathcal{F}(\mathcal{V}_j);$
- 5: **for** i = 1, ..., j **do**
- 6:  $h_{i,j} = \mathcal{U}_{p_{i,1},p_{i,2},...,p_{i,N}};$
- 7:  $\mathcal{U} = \mathcal{U} \dot{h}_{i,j} \dot{\mathcal{V}}_i;$
- 8: end for
- 9: Determine triple  $(i_{1,0}, i_{2,0}, \ldots, i_{N,0})$  using Algorithm 2 for the input  $\mathcal{U}$ ;
- 10: Set  $h_{j+1,j} = \mathcal{U}_{i_{1,0},i_{2,0},\dots,i_{N,0}}$ ;  $\mathcal{V}_{j+1} = \mathcal{U}/h_{j+1,j}$ ;  $p_{j+1,\eta} = i_{\eta,0}$  for  $\eta = 1, 2, \dots, N$ ;

11: end for