# Determining the Distribution Route Using the Clustered Generalised Vehicle Routing Problem Model and Dijkstra's Algorithm 

Awal Aflizal Zubir ${ }^{1 *}$, Andriansyah ${ }^{1}$, Salsabila Derisma ${ }^{1}$<br>Industrial Engineering Program, Mechanical and Industrial Engineering Department, Universitas Syiah Kuala, JI. Tgk Abdur<br>Rauf No.7, Banda Aceh 23111, Indonesia<br>* Corresponding email: awalaflizal.zubir@unsyiah.ac.id

## Riwayat Artikel <br> Diterima <br> 01/08/2022 <br> Disetujui <br> 11/08/2022 <br> Diterbitkan <br> 30/09/2022


#### Abstract

Abstrak Kegiatan pada sektor industri tidak terpisahkan dari proses distribusi dan logistik, salah satu sektor yang erat kaitannya dengan distribusi adalah penyaluran LPG 3 kg . Permasalahan paling umum yang terjadi pada proses distribusi adalah penentuan rute kendaraan yang sering disebut dengan Vehicle Routing Problem. PT Amalia Yusri sebagai salah satu agen penyalur LPG 3 kg di Banda Aceh dan tidak terlepas dari permasalahan penentuan rute distribusi. Salah satu alternatif penentuan rute distribusi adalah menggunakan model Clustered Generalized Vehicle Routing Problem yang bertujuan untuk mencari rute optimal pada titik tujuan yang telah dikelompokan. Berdasarkan jumlah permintaan dan kapasitas kendaraan, maka cluster dikelompokkan dan dibentuk menjadi dua hari yang terbagi ke dalam lima kelompok. Hari pertama terdiri dari kelompok 1 dan kelompok 2, sedangkan hari kedua terdiri dari kelompok 3, kelompok 4, dan kelompok 5. Metode yang akan diuji untuk mengoptimalkan jarak pendistribusian adalah Algoritma Dijkstra. Dari penerapan model CGVRP dan Algoritma Dijkstra diperoleh total jarak untuk kelompok 1 sebesar 25.582 m , total jarak untuk kelompok 2 sebesar 24.650 m , total jarak untuk kelompok 3 sebesar 39.350 m , total jarak untuk kelompok 4 sebesar 27.500 m , dan total jarak untuk kelompok 5 sebesar 38.500 m .


Kata Kunci: Rute Distribusi Logistik, LPG, CGVRP, Algoritma Dijkstra


#### Abstract

Distributions and logistics processes are the essential element activities integral to the industrial sector, one of them is the distribution of LPG 3 kg . The most common issue that occurs in the distribution process is determining the vehicle routes which are often called Vehicle Routing Problems. PT Amalia Yusri is one of the distributor agents of LPG 3 kg in Banda Aceh and also had the same issue to determine the distribution routes. The Clustered Generalized Vehicle Routing Problem model is used to determine the optimal distribution route on the vertex destinations that have been grouped. Based on the number of requests and vehicle capacity, the grouping and formation of clusters are divided into two days and split into five groups. The first day consists of group 1 and group 2; while on the second day are group 3, group 4, and group 5. The method that is used to optimise distribution distance is Dijkstra Algorithm. From the application of the CGVRP model and Dijkstra Algorithm obtained a total distance for group 1 is 25.58 Km , group 2 is 24.65 Km , group 3 is $39,35 \mathrm{Km}$, group 4 is 27.5 km , and group 5 is 38.5 km .


Keywords: Logistic Distribution Route, LPG, CGVRP, Dijkstra Algorithm

## 1. Introduction

The logistics process is directly connected to distribution process activities, specifically in the industrial sector where the process of supply from production to consumers is affected by the change of optimisation of costs, product quality and delivery time [1]-[3]. The standpoint of logistics is affiliated with planning, distribution, and the transportation linked with factories to warehouses and then to the market location [2], [4]. Technically speaking, the distribution of goods is carried out by the distributor that used vehicles which has different sizes of capacity [5].

The issues of distribution routes are often referred to as Vehicle Routing Problems (VRP) which are related to mathematically analysis to determine vehicle distribution routes from the distribution centre to customers located in different locations [6]. One variation of VRP is the Vehicle Routing Problem with Simultaneous Deliveries and Pickups (VRPSDP) where vehicles are not only necessary to deliver goods to the customer's location but also the delivery and pick-up of goods carried out simultaneously in one vehicle stop [7]. In order to, solve the problems that occur in the distribution process of 3 kg LPG at PT Amalia Yusri in Banda Aceh, this paper uses the Clustered Generalized Vehicle Routing Problem (CGVRP) modelling and Dijkstra's Algorithm.

The CGVRP model is a development of the Generalized Vehicle Routing Problem (GVRP) method in which the distribution is clustered or grouped based on predetermined criteria [8]. The method that will be tested to optimize the distribution distance is Dijkstra's Algorithm, where the algorithm is used to solve the problem of finding the optimal route by calculating the smallest weight of each point [9].

According to the information obtained from PT Amalia Yusri, the current distribution process of Liquefied Petroleum Gas (LPG) cylinders carried out by PT Amalia Yusri has only relied on the intuition and experience of truck drivers. This is of course less effective and efficient because the driver does not analyse in determine the distribution route, so it is not known whether the route currently used is the optimal route [8]. In order to determine the finest solution for PT Amalia Yusri, the implementation of the CGVRP model and Dijkstra's Algorithm is conducted to present the optimal alternative solutions.

## 2. Research Methods

### 2.1. Data Modelling

CGVRP is a variation of GVRP as a general issue of vehicle route clusters, where all vertices in each cluster must be visited sequentially in the vehicle route. CGVRP can determine the minimum cost collection of $m$ vehicle routes exactly once by performing a Hamilton path at each vertex, without
exceeding the carrying capacity of the vehicle [10], [11].


Figure 1. Example of CGVRP Solutions [9], [12]
Figure 1 shows an example of a feasible solution for CGVRP of a distribution process of goods with the vehicle capacity of $Q=25$ cylinders, the distribution process is executed by using two identical vehicles. The first vehicle starts distribution from the warehouse $\left(v_{0}\right)$ to vertex $v_{3}$ with a request for $d_{3}=4$ cylinders then goes to vertex $v_{2}$ with $d_{2}=5$ cylinders, and then goes to vertex $v_{1}$ with $d_{1}=3$ cylinders; These vertices are vertex in cluster $1\left(V_{1}\right)$ with a total number of requests $q_{1}=12$ cylinders. Then the first vehicle goes to vertex $v_{4}$ with the number of requests $d_{4}=5$ cylinders, and continues to vertex $v_{5}$ with the number of requests $d_{5}$ $=4$ cylinders; vertices $v_{4}$ and $v_{5}$ are vertices in cluster $2\left(V_{2}\right)$ with the number of requests $q_{2}=9$ cylinders. Then the first vehicle returns to $v o$. The same process occurred in the second vehicle assigned to cluster 5, cluster 4 and cluster 3 [12]. The optimal route of the distribution process is the total distance travelled by the two vehicles.

### 2.2. Dijkstra Algorithm

Dijkstra's algorithm is one of the variants of the Greedy algorithm that is included in one of the popular algorithms to solve problems related to optimisation. To determine the solution, Dijkstra's algorithm uses the Greedy principle, to find the optimal solution at each step to get the best solution. Dijkstra's algorithm works by using the priority queue principle, in which the point that has the highest priority will be traced. In determining the priority points, Dijkstra's algorithm compares each value or weight of the points that are at one level. Furthermore, the value of the point is saved to be compared with the value that will be discovered from the route that was just found thereupon, and so on until the point needed for inspection is found [13], [14].

Dijkstra's algorithm requires the parameters of the place of origin and destination, where the final result
of the algorithm is the shortest distance from the place of origin to the destination and its route [15]. The procedural steps of Dijkstra's algorithm are as follows [12], [14], [16]-[19]:

1. Initialisation: $L=\{ \} ; V=\left\{v_{1}, v_{2}, \ldots, v_{n}\right\}$.
2. For $i=2,3, \ldots, \mathrm{n}$, do $D(i)=W(1, j)$.
3. As long as $v_{n} \notin L$ ( $v_{n}$ is not a permanent point), do: a. Choose $v_{n} \in V-L$ (non-permanent vertex) with the smallest $D(k) . L=L \cup\left\{v_{k}\right\}$ (make $v_{k}$ as the permanent point).
b. For every $v_{j} \in V-L$ do: if $D(k)+W(k, j)<D$ $(j)$ then replace $D(j)$ with $D(k)+W(k, j)$.

## 3. Results and Analysis

### 3.1 Grouping and Cluster Formation

PT Amalia Yusri distributes LPG cylinders of 3 kg every day to 50 gas stations in Banda Aceh City using two vehicles, that have the same capacity. The grouping is classified by taking into matter the vehicle capacity with the total number of requests not exceeding the maximum vehicle capacity of 260 of 3 kg LPG cylinders. If the number of requests exceeds the vehicle capacity, a new group will be formed.

Based on the number of requests and vehicle capacity, the clusters are grouped and formed into two days which are divided into five groups. The first day consisted of group 1 and group 2, while the second day consisted of group 3, group 4, and group 5, with
one group consisting of 2 trucks operating. Hence, by using Dijkstra Algorithm that determining the shortest travel distance at a one-time process of delivering process can be concluded as the optimum routes. Therefore, the results will serve as a guideline for the drivers and company while distributing the goods of LPG 3 Kg . Through the route order of delivering one trip route determined of both vehicles need to execute, it can save the cost of time and fuel consumption.

### 3.2. Distance Determination

Determination of the distance from the warehouse to each LPG base and the distance between each base in the cluster is utilised with the Google Maps application.


Figure 2. Route measurement using the Google Maps application

Table 1. LPG station grouping

| Sub-district | Station name | Store code | Request quantity | Truck capacity | Truck |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Punge Ujong | Berkah Gas II | 13 | 12 | 248 | 1 |
|  | Moeda Bejaya | 27 | 12 |  |  |
|  | Punge Gas | 33 | 12 |  |  |
|  | UD. Puma Kelana | 43 | 12 |  |  |
| Punge Jurong | Anggrek | 9 | 29 |  |  |
|  | Umaira | 46 | 29 |  |  |
|  | Usaha Loen | 49 | 29 |  |  |
| Gampong Baro | Afraff | 5 | 29 |  |  |
| Lambung | Tekad Gas | 34 | 21 |  |  |
| Deah Baro | CV. Mama Energi | 15 | 17 |  |  |
| Surien | Dersi Gas | 50 | 30 |  |  |
| Asoe Nanggroe | Amiruddin Gas | 7 | 17 |  |  |
| Punge Blang Cut | Berkah Gas | 12 | 135 | 236 | 2 |
| Bitai | CV. Pratama Cipta Jaya | 17 | 29 |  |  |
| Lampoh Daya | Nuansa Barokah | 30 | 35 |  |  |
| Lamjamee | UD. Khadafi | 41 | 37 |  |  |
| Emperom | Bunda Gas | 14 | 30 | 223 | 1 |
|  | Usaha Baru | 47 | 30 |  |  |
| Lamteumen Barat | UD. Mulia Gas | 42 | 59 |  |  |
| Geuceu Meunara | UD. Saniah | 44 | 56 |  |  |
| Lampaseh Kota | Dedex Gas | 18 | 48 |  |  |
| Gampong Jawa | Andri Gas | 8 | 49 | 165 | 2 |
| Gampong Pande | Adia Ari Gas | 4 | 17 |  |  |


| Sub-district | Station name | Store code | Request quantity | Truck capacity | Truck |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Sukaramai | 14.232.485 | 3 | 99 |  |  |
| Peuniti | Maju Jaya | 24 | 64 | 224 | 1 |
|  | UD. Hawana | 40 | 64 |  |  |
| Ateuk Jawo | Barokah Gas | 11 | 21 |  |  |
|  | Nasriah Johan | 29 | 21 |  |  |
| Lam Ara | UD. Cek Poel | 38 | 54 |  |  |
| Mibo | Agaz Water | 6 | 12 | 236 | 2 |
|  | CV. Mita Bacut | 16 | 12 |  |  |
|  | UD. Fahrol Gas | 39 | 12 |  |  |
| Panteriek | Toko Laris | 35 | 89 |  |  |
| Lamseupeung | Usaha Baru Gas | 48 | 41 |  |  |
| Cot Mesjid | Pintoe Raseuki | 32 | 69 |  |  |
| Lamdom | Pinto Aceh | 31 | 34 | 252 | 1 |
| Laksana | Kana Usaha Gas | 21 | 112 |  |  |
| Mulia | 14.231.482 | 2 | 106 |  |  |
| Keuramat | H Amin Gas | 19 | 55 | 215 | 2 |
|  | Tsani Gas | 36 | 55 |  |  |
| Beurawe | Aulia Gas | 10 | 53 |  |  |
|  | Kedai Hemi | 22 | 53 |  |  |
| Lamdingin | Mulia Sakti | 28 | 69 | 182 | 1 |
| Lambaro Skep | Keupula | 23 | 57 |  |  |
|  | Makmur Gas | 25 | 57 |  |  |
| Kota Baru | Menjemput Rezeki | 26 | 27 | 257 | 2 |
| le Masen Ulee Kareng | UD. Tabarak | 45 | 45 |  |  |
| Peurada | UB. Al-Mukmin | 37 | 69 |  |  |
| Lamgugob | 14.231 .450 | 1 | 58 |  |  |
|  | Idola Remaja | 20 | 58 |  |  |

### 3.3 Dijkstra Algorithm application constructing CGVRP model

The beginning of destination vertex for each cluster and the analysing calculations using the Dijkstra algorithm to develop a CGVRP model for each group were done with the Google Maps application. Therefore, the distance between warehouses to each

LPG station and the distance between the LPG station assigned to each vehicle can be determined. In group 1 the first vehicle was assigned to visit seven clusters. Determination of the initial destination vertex of the first vehicle, specifically the point that has the closest distance from the seven clusters.

Table 2. Distance matrix first vehicle of group 1 ( m )

| Stations code | $\mathbf{0}$ | $\mathbf{1 3}$ | $\mathbf{2 7}$ | $\mathbf{3 3}$ | $\mathbf{4 3}$ | $\mathbf{9}$ | $\mathbf{4 6}$ | $\mathbf{4 9}$ | $\mathbf{5}$ | $\mathbf{3 4}$ | $\mathbf{1 5}$ | $\mathbf{5 0}$ | $\mathbf{7}$ |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 0 |  |  |  |  |  |  |  |  |  |  |  |  |
| 13 | 190 | 0 |  |  |  |  |  |  |  |  |  |  |  |
| 27 | 200 | 170 | 0 |  |  |  |  |  |  |  |  |  |  |
| 33 | 220 | 190 | 100 | 0 |  |  |  |  |  |  |  |  |  |
| 43 | 210 | 64 | 230 | 220 | 0 |  |  |  |  |  |  |  |  |
| 9 | 850 | 800 | 700 | 700 | 900 | 0 |  |  |  |  |  |  |  |
| 46 | 900 | 800 | 650 | 600 | 850 | 280 | 0 |  |  |  |  |  |  |
| 49 | 950 | 900 | 750 | 650 | 1000 | 92 | 350 | 0 |  |  |  |  |  |
| 5 | 2100 | 1800 | 2000 | 1900 | 1700 | 2400 | 2700 | 2400 | 0 |  |  |  |  |
| 34 | 1800 | 1100 | 1300 | 1300 | 1000 | 1700 | 2000 | 1800 | 1800 | 0 |  |  |  |
| 15 | 1600 | 1700 | 2000 | 1900 | 1600 | 2300 | 2500 | 2400 | 2400 | 1200 | 0 |  |  |
| 50 | 3300 | 3100 | 3400 | 3400 | 3100 | 3600 | 4000 | 3500 | 1700 | 2600 | 3800 | 0 |  |
| 7 | 3700 | 3700 | 3900 | 3900 | 3000 | 3700 | 4000 | 3800 | 2200 | 2600 | 2800 | 1400 | 0 |

Based on the distance matrix from table 2, it shows the initial destination vertex is the base with code 13 since it has the closest distance from the starting point, which PT Amalia Yusri has a code of 0.

Additionally, calculations are conducted using Dijkstra's Algorithm, the results of the analysis of these calculations are shown in table 3.

Table 3. Application of Dijkstra's Algorithm to the First Vehicle Distribution of Group 1

| Iteration | 0 | 13 | 27 | 33 | 43 | 9 | 46 | 49 | 5 | 34 | 15 | 50 | 7 | L |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 0 | $\infty$ | $\infty$ | $\infty$ | $\infty$ | $\infty$ | $\infty$ | $\infty$ | $\infty$ | $\infty$ | $\infty$ | $\infty$ | $\infty$ | \{0\} |
| 2 | - | $\begin{gathered} \operatorname{Min} \\ \{\infty, 0 \\ +190\} \\ =190 \\ \hline \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{\infty, 0 \\ +200\} \\ =200 \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{\infty, 0 \\ +220\} \\ =220 \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{\infty, 0 \\ +210\} \\ =210 \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{\infty, 0 \\ +850\} \\ =850 \\ =85 \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{\infty, 0 \\ +900\} \\ =900 \\ \hline \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{\infty, 0 \\ +950\} \\ =950 \\ =95 \end{gathered}$ | $\begin{gathered} \hline \operatorname{Min} \\ \{\infty, 0 \\ +2100\} \\ =2100 \\ \hline \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{\infty, 0 \\ +1800\} \\ =1800 \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{\infty, 0 \\ +1600\} \\ =1600 \\ \hline \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{\infty, 0 \\ +3300\} \\ =3300 \\ \hline \end{gathered}$ | $\begin{gathered} \text { Min } \\ \{\infty, 0 \\ +3700\} \\ =3700 \end{gathered}$ | $\{0,13\}$ |
| 3 | - | - | $\begin{gathered} \operatorname{Min} \\ \{200, \\ 190 \\ +170\} \\ =200 \\ \hline \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{220, \\ 190 \\ +190\} \\ =220 \\ \hline \end{gathered}$ | $\begin{gathered} \text { Min } \\ \{210, \\ 190 \\ +64\} \\ =210 \\ \hline \end{gathered}$ | $\begin{gathered} \mathrm{Min} \\ \{850, \\ 190 \\ +800\} \\ =850 \\ \hline \end{gathered}$ | $\begin{gathered} \mathrm{Min} \\ \{900, \\ 190 \\ +800\} \\ =900 \\ \hline \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{950, \\ 190 \\ +900\} \\ =950 \\ \hline \end{gathered}$ | $\begin{gathered} \hline \operatorname{Min} \\ \{2100, \\ 190 \\ +1800\} \\ =1990 \\ \hline \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{1800, \\ 190 \\ +1100\} \\ =1290 \\ \hline \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{1600, \\ 190 \\ +1700\} \\ =1600 \\ \hline \end{gathered}$ | $\begin{gathered} \mathrm{Min} \\ \{3300, \\ 190 \\ +3100\} \\ =3290 \\ \hline \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{3700, \\ 190 \\ +3700\} \\ =3700 \\ \hline \end{gathered}$ | $\begin{gathered} \{0,13, \\ 27\} \end{gathered}$ |
| 4 | - | - | - | $\begin{gathered} \mathrm{Min} \\ \{220, \\ 200 \\ +100\} \\ =220 \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{210, \\ 200 \\ +230\} \\ =210 \end{gathered}$ | $\begin{gathered} \text { Min } \\ \{850, \\ 200 \\ +700\} \\ =850 \end{gathered}$ | $\begin{gathered} \hline \text { Min } \\ \{900, \\ 200 \\ +650\} \\ =850 \end{gathered}$ | $\begin{gathered} \text { Min } \\ \{950, \\ 200 \\ +750\} \\ =950 \\ \hline \end{gathered}$ | $\begin{gathered} \hline \text { Min } \\ \{1990, \\ 200 \\ +2000\} \\ =1990 \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{1290, \\ 200 \\ +1300\} \\ =1290 \end{gathered}$ | $\begin{gathered} \text { Min } \\ \{1600, \\ 200 \\ +2000\} \\ =1600 \end{gathered}$ | $\begin{gathered} \text { Min } \\ \{3290, \\ 200 \\ +3400\} \\ =3290 \end{gathered}$ | $\begin{gathered} \text { Min } \\ \{3700, \\ 200 \\ +3900\} \\ =3700 \end{gathered}$ | $\begin{aligned} & \{0,13, \\ & 27,43\} \end{aligned}$ |
| 5 | - | - | - | $\begin{gathered} \operatorname{Min} \\ \{220, \\ 210 \\ +220\} \\ =220 \\ \hline \end{gathered}$ | - | $\begin{gathered} \mathrm{Min} \\ \{850, \\ 210 \\ +900\} \\ =850 \\ \hline \end{gathered}$ | $\begin{gathered} \hline \operatorname{Min} \\ \{850, \\ 210 \\ +850\} \\ =850 \\ \hline \end{gathered}$ | $\begin{gathered} \text { Min } \\ \{950, \\ 210 \\ +1000\} \\ =950 \\ \hline \end{gathered}$ | $\begin{gathered} \hline \operatorname{Min} \\ \{1990, \\ 210 \\ +1700\} \\ =1910 \\ \hline \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{1290, \\ 210 \\ +1000\} \\ =1210 \\ \hline \end{gathered}$ | $\begin{gathered} \text { Min } \\ \{1600, \\ 210 \\ +1600\} \\ =1600 \\ \hline \end{gathered}$ | $\begin{gathered} \hline \text { Min } \\ \{3290, \\ 210 \\ +3100\} \\ =3290 \\ \hline \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{3700, \\ 210 \\ +3000\} \\ =3210 \\ \hline \end{gathered}$ | $\begin{gathered} \{0,13, \\ 27,43, \\ 33\} \end{gathered}$ |
| 6 | - | - | - | - | - | $\begin{gathered} \text { Min } \\ \{850, \\ 220 \\ +700\} \\ =850 \end{gathered}$ | $\begin{gathered} \hline \operatorname{Min} \\ \{850, \\ 220 \\ +600\} \\ =820 \end{gathered}$ | $\begin{gathered} \text { Min } \\ \{950, \\ 220 \\ +650\} \\ =870 \end{gathered}$ | $\begin{gathered} \hline \text { Min } \\ \{1910, \\ 220 \\ +1900\} \\ =1910 \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{1210, \\ 220 \\ +1300\} \\ =1210 \end{gathered}$ | $\begin{gathered} \text { Min } \\ \{1600, \\ 220 \\ +1900\} \\ =1600 \end{gathered}$ | $\begin{gathered} \hline \operatorname{Min} \\ \{3290, \\ 220 \\ +3400\} \\ =3290 \\ \hline \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{3210, \\ 220 \\ +3900\} \\ =3210 \end{gathered}$ | $\begin{aligned} & \{0,13, \\ & 27,43, \\ & 33,46\} \end{aligned}$ |
| 7 | - | - | - | - | - | $\begin{gathered} \text { Min } \\ \{850, \\ 820 \\ +280\} \\ =850 \end{gathered}$ | - | $\begin{gathered} \operatorname{Min} \\ \{870, \\ 820 \\ +350\} \\ =870 \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{1910, \\ 820 \\ +2700\} \\ =1910 \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{1210, \\ 820 \\ +2000\} \\ =1210 \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{1600, \\ 820 \\ +2500\} \\ =1600 \end{gathered}$ | $\begin{gathered} \mathrm{Min} \\ \{3290, \\ 820 \\ +4000\} \\ =3290 \\ \hline \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{3210, \\ 820 \\ +4000\} \\ =3210 \end{gathered}$ | $\begin{gathered} \{0,13, \\ 27,43, \\ 33,46, \\ 9\} \end{gathered}$ |
| 8 | - | - | - | - | - | - | - | $\begin{gathered} \operatorname{Min} \\ \{870, \\ 850 \\ +92\}= \\ 870 \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{1910, \\ 850 \\ +2400\} \\ =1910 \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{1210, \\ 850 \\ +1700\} \\ =1210 \\ \hline \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{1600, \\ 850 \\ +2300\} \\ =1600 \end{gathered}$ | $\begin{gathered} \mathrm{Min} \\ \{3290, \\ 850 \\ +3600\} \\ =3290 \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{3210, \\ 850 \\ +3700\} \\ =3210 \end{gathered}$ | $\begin{aligned} & \{0,13, \\ & 27,43, \\ & 33,46, \\ & 9,49\} \end{aligned}$ |
| 9 | - | - | - | - | - | - | - | - | $\begin{gathered} \hline \operatorname{Min} \\ \{1910, \\ 870 \\ +2400\} \\ =1910 \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{1210, \\ 870 \\ +1800\} \\ =1210 \end{gathered}$ | $\begin{gathered} \text { Min } \\ \{1600, \\ 870 \\ +2400\} \\ =1600 \end{gathered}$ | $\begin{gathered} \hline \operatorname{Min} \\ \{3290, \\ 870 \\ +3500\} \\ =3290 \\ \hline \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{3210, \\ 870 \\ +3800\} \\ =3210 \end{gathered}$ | $\begin{gathered} \{0,13, \\ 27,43, \\ 33,46, \\ 9,49, \\ 34\} \\ \hline \end{gathered}$ |
| 10 | - | - | - | -- | - | - | - | - | $\begin{gathered} \operatorname{Min} \\ \{1910, \\ 1210 \\ +1800\} \\ =1910 \end{gathered}$ | - | $\begin{gathered} \operatorname{Min} \\ \{1600, \\ 1210 \\ +1200\} \\ =1600 \end{gathered}$ | $\begin{gathered} \mathrm{Min} \\ \{3290, \\ 1210 \\ +2600\} \\ =3290 \\ \hline \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{3210, \\ 1210 \\ +2600\} \\ =3210 \end{gathered}$ | $\begin{gathered} \{0,13, \\ 27,43, \\ 33,46, \\ 9,49, \\ 34,15\} \end{gathered}$ |
| 11 | - | - | - | - | - | - | - | - | $\begin{gathered} \operatorname{Min} \\ \{1910, \\ 1600 \\ +2400\} \\ =1910 \end{gathered}$ | - | - | $\begin{gathered} \operatorname{Min} \\ \{3290, \\ 1600 \\ +3800\} \\ =3290 \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{3210, \\ 1600 \\ +2800\} \\ =3210 \end{gathered}$ | $\begin{gathered} \{0,13, \\ 27,43, \\ 33,46, \\ 9,49, \\ 34,15, \\ 5\} \end{gathered}$ |
| 12 | - | - | - | - | - | - | - | - | - | - | - | $\begin{gathered} \operatorname{Min} \\ \{3290, \\ 1910 \\ +1700\} \\ =3290 \end{gathered}$ | $\begin{gathered} \operatorname{Min} \\ \{3210, \\ 1900 \\ +2200\} \\ =3210 \end{gathered}$ | $\begin{gathered} \{0,13, \\ 27,43, \\ 33,46, \\ 9,49, \\ 34, \\ 15,5,7\} \end{gathered}$ |
| 13 | - | - | - | - | - | - | - | - | - | - | - | $\begin{gathered} \operatorname{Min} \\ \{3290, \\ 3210 \\ +1400\} \\ =3290 \end{gathered}$ | - | $\begin{gathered} \{0,13, \\ 27,43, \\ 33,46, \\ 9,49, \\ 34,15, \\ 5,7, \\ 50\} \end{gathered}$ |

Based on the results of the implementation of the Dijkstra Algorithm listed in table 3, the distribution route for the first vehicle in group 1 starts from the warehouse (0) to vertex $13 \rightarrow$ vertex $27 \rightarrow$ vertex 43 $\rightarrow$ vertex $33 \rightarrow$ vertex $46 \rightarrow$ vertex $9 \rightarrow$ vertex $49 \rightarrow$ vertex $34 \rightarrow$ vertex $15 \rightarrow$ vertex $5 \rightarrow$ vertex $7 \rightarrow$ vertex 50. The same way applied to the second vehicle and attained the distribution route starting from the warehouse ( 0 ) to vertex $12 \rightarrow$ vertex $17 \rightarrow$ vertex 41 $\rightarrow$ vertex 30. In addition, the CGVRP modelling distribution route is shown in Figure 3.


Figure 3. CGVRP model route distribution of group 1
The same steps were applied to the other groups and obtained a CGVRP model for all groups as shown in the following figures.


Figure 4. CGVRP model route distribution of group 2


Figure 5. CGVRP model route distribution of group 3


Figure 6. CGVRP model route distribution of group 4


Figure 7. CGVRP model route distribution of group 5
The calculation of the mileage of groups 1 to 5 is based on table 3 on determining the distribution route of 3 kg LPG by PT Amalia Yusri using the Dijkstra algorithm shown in table 4.

Table 4. Calculation of the proposed route distance

| Group | Vehicle | Route | Vehicle capacity | Total Capacity | Travel distance (m) | Total Travel (m) |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | Vehicle 1 | PT Amalia Yusri (0) $\rightarrow$ Berkah Gas II (13) $\rightarrow$ Moeda Bejaya (27) $\rightarrow$ UD. Puma Kelana (43) $\rightarrow$ Punge Gas (33) $\rightarrow$ Umaira (46) $\rightarrow$ Anggrek (9) $\rightarrow$ Usaha Loen (49) $\rightarrow$ Tekad Gas (34) $\rightarrow$ CV. Mama Energi (15) $\rightarrow$ Afraff (5) $\rightarrow$ Amiruddin Gas (7) $\rightarrow$ Dersi Gas $(50) \rightarrow$ PT Amalia Yusri (0) | 248 | 484 | $\begin{aligned} & 190+170+230 \\ & +220+600+ \\ & 280+92+1800 \\ & +1200+2400+ \\ & 2200+1400+ \\ & 3300=14082 \end{aligned}$ | 25,582 |


| Group | Vehicle | Route | Vehicle capacity | Total Capacity | Travel distance (m) | Total Travel (m) |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Vehicle 2 | PT Amalia Yusri (0) $\rightarrow$ Berkah Gas (12) $\rightarrow$ CV. Pratama Cipta Jaya (17) $\rightarrow$ UD. Khadafi (41) $\rightarrow$ Nuansa Barokah (30) $\rightarrow$ PT Amalia Yusri (0) | 236 |  | $\begin{aligned} & 2000+1100+ \\ & 2000+1600+ \\ & 4800=11500 \end{aligned}$ |  |
| 2 | Vehicle 1 | PT Amalia Yusri (0) $\rightarrow$ Dedex Gas (18) $\rightarrow$ Bunda Gas (14) $\rightarrow$ Usaha Baru (47) $\rightarrow$ UD. Mulia Gas (42) $\rightarrow$ UD Saniah (44) $\rightarrow$ PT Amalia Yusri (0) | 223 | 388 | $\begin{aligned} & 1700+4800+ \\ & 600+1100+ \\ & 1900+4800= \\ & 14900 \end{aligned}$ | 24,650 |
|  | Vehicle 2 | PT Amalia Yusri (0) $\rightarrow 14.232 .485$ <br> (3) $\rightarrow$ Adia Ari Gas (4) $\rightarrow$ Andria Gas <br> (8) $\rightarrow$ PT Amalia Yusri (0) | 165 |  | $\begin{aligned} & 2200+3300+ \\ & 750+3500= \\ & 9750 \end{aligned}$ |  |
| 3 | Vehicle 1 | PT Amalia Yusri (0) $\rightarrow$ Maju Jaya (24) $\rightarrow$ UD. Hawana (40) $\rightarrow$ Barokah Gas (11) $\rightarrow$ Nasriah Johan (29) $\rightarrow$ UD. Cek Poel (38) $\rightarrow$ PT Amalia Yusri (0) | 224 | 460 | $\begin{aligned} & 3500+450+ \\ & 1900+400+ \\ & 4000+6600= \\ & 16850 \end{aligned}$ | 39,350 |
|  | Vehicle 2 | PT Amalia Yusri (0) $\rightarrow$ UD. Fahrol Gas (39) $\rightarrow$ Agaz Water (6) $\rightarrow$ CV. Mita Bacut (16) $\rightarrow$ Usaha Baru Gas (48) $\rightarrow$ Pintoe Raseuki (32) $\rightarrow$ Toko Laris (35) $\rightarrow$ PT Amalia Yusri (0) | 236 |  | $\begin{aligned} & 3700+1000+ \\ & 800+5000+ \\ & 3300+3100+ \\ & 5600=22500 \end{aligned}$ |  |
| 4 | Vehicle 1 | PT Amalia Yusri (0) $\rightarrow$ Kana Usaha Gas (21) $\rightarrow$ 14.231.482 (2) $\rightarrow$ Pinto Aceh (31) $\rightarrow$ PT Amalia Yusri (0) | 252 | 467 | $\begin{aligned} & 3700+450+ \\ & 4700+6500= \\ & 15350 \end{aligned}$ | 27,500 |
|  | Vehicle 2 | PT Amalia Yusri (0) $\rightarrow \mathrm{H}$. Amin Gas $(19) \rightarrow$ Tsani Gas (36) $\rightarrow$ Kedai Helmi (22) $\rightarrow$ Aulia Gas (10) $\rightarrow \rightarrow$ PT Amalia Yusri (0) | 215 |  | $\begin{aligned} & 3300+600+ \\ & 1400+650+ \\ & 6200=12150 \end{aligned}$ |  |
| 5 | Vehicle 1 | PT Amalia Yusri (0) $\rightarrow$ Mulia Sakti (28) $\rightarrow$ Keupula (23) $\rightarrow$ Makmur Gas (25) $\rightarrow$ PT Amalia Yusri (0) | 182 | 439 | $\begin{aligned} & 5100+1500+ \\ & 2100+6000= \\ & 14700 \end{aligned}$ | 38,500 |
|  | Vehicle 2 | PT Amalia Yusri (0) $\rightarrow$ Menjemput Rezeki (26) $\rightarrow$ UD. Tabarak (45) $\rightarrow$ UB. Al-Mukmin (37) $\rightarrow$ 14.231.450 (1) $\rightarrow$ Idola Remaja (20) $\rightarrow$ PT Amalia Yusri (0) | 257 |  | $\begin{aligned} & 6700+4000+ \\ & 2700+1600+ \\ & 500+8300= \\ & 23800 \end{aligned}$ |  |

### 3.4 Fuel Consumption Cost Based on CGVRP Model Route

The fuel consumption of the total travel distance of both vehicles will be calculated to determine the optimum route with the assumption that the vehicle used had the same fuel capacity ratio which 9 $\mathrm{Km} / \mathrm{litre}$ using Pertalite gasoline 7,650 IDR as the vehicle fuel without considering traffic light, vehicle full load capacity weight. Then, from table 4 the rate of fuel consumption is as follows; the sum of vehicle 1 travel is 75.8 km with fuel consumption of 8.43 litre that cost 64,490 IDR. Sum vehicle 2 travel is 79.7 km with fuel consumption of 8.86 litres that cost 67,779 IDR. Thus, the average fuel cost for two days delivering $3 K G$ LPG is 66,135 IDR.

## 4. Conclusion

The route obtained from this research results of the CGVRP model and the Dijkstra Algorithm will serve as a guideline for the drivers while delivering
the goods of PT Amalia Yusri which LPG 3 Kg . Based on the number of requests and vehicle capacity, the grouping and clustering of the CGVRP model and the Dijkstra Algorithm are divided into two days and then divided more into five groups with the results of calculating the total distance of 25.58 km (group 1), 24.65 km (group 2), 39.35 km (group 3), 27.5 km (group 4), and 38.5 km (group 5). It is concluded that vehicles 1 and 2 consumed 8.43 and 8.86 litres of Pertalite gasoline that cost around 66,135 IDR.

## References

[1] D. W. Nugraha, A. Y. E. Dodu, and S. Septiana, "Sistem Penentuan Rute Pendistribusian Produk Air Mineral Menggunakan Algoritma Ant Colony System," ILKOM Jurnal Ilmiah, vol. 11, no. 2, pp. 86-94, 2019, doi: 10.33096/ilkom.v11i2.418.86-94.
[2] F. R. Zanjirani, S. Rezapour, and L. Kardar, Logistics Operations and Management. London: Elsevier, 2011. doi: 10.1016/B978-0-12-385202-1.X0001-1.
[3] Alan. Rushton, P. Croucher, and P. Baker, The handbook of logistics and distribution management: understanding the supply chain, 5th ed. London: Kogan Page Limited, 2014.
[4] A. Sutoni and N. Asilah, "Penentuan Jarak Pendistribusian Gas LPG 3 Kg Dengan Metode Algoritma Heuristik," Jurnal Media Teknik dan Sistem Industri, vol. 2, no. 2, pp. 37-42, 2018, doi: 10.35194/jmtsi.v2i2.416.
[5] S. D. Juniarto, E. M. K, A. Fariza, and I. Prasetyaningrum, "Optimasi Distribusi Barang Berdasarkan Rute Dan Daya Tampung Menggunakan Metode Simulated Annealing," Jurnal Teknik Informatika, PENS - ITS Surabaya, 2011.
[6] F. S. Lubis and M. K. Herliansyah, "Vehicle Routing Problem with Simultaneous Delivery and Pick-up Services (VRPSDP) pada Distribusi Tabung Gas LPG 3 Kg (Kasus: PT. Lentera Putera Sejahtera)," Seminar Nasional Teknik Industri, pp. 22-29, 2017.
[7] A. S. Tasan and M. Gen, "A genetic algorithm based approach to vehicle routing problem with simultaneous pick-up and deliveries," Computers and Industrial Engineering, vol. 62, no. 3, pp. 755-761, 2012, doi: 10.1016/j.cie.2011.11.025.
[8] K. Hermanto, I. Adiasa, S. Altarisi, R. Rabani, and M. Amirul, "Rute Usulan Pendistribusian LPG Menggunakan Model Clustered Generalized Vehicle Routing Problem (CGVRP) dan Algoritma Dijkstra," Performa: Media Ilmiah Teknik Industri, vol. 19, no. 1, pp. 27-36, 2020, doi: 10.20961/performa.19.1.41858.
[9] K. Hermanto and T. D. Ermayanti, "Analisa Optimasi Rute Transportasi Antar Jemput

Siswa Menggunakan Model CGVRP dan Algoritma Dijkstra di SDIT Darus Sunnah," Unisda Journal of Mathematics and Computer Science (UJMC), vol. 5, no. 2, pp. 19-28, 2019, doi: 10.52166/ujmc.v5i2.1653.
[10] K. Hermanto, E. Ruskartina, and Gunawan, "Application of the CGVRP model and the 01 linear programming determines the RASKIN (subsidized rice) distribution route," Journal of Physics: Conference Series, vol. 1778, no. 1, pp. 1-8, 2021, doi: 10.1088/1742-6596/1778/1/012002.
[11] P. C. Pop, Generalized Network Design Problems: Modeling and Optimization. De Gruyter, 2012. doi:10.1515/9783110267686.
[12] K. Hermanto and E. Ruskartina, "Optimasi Rute Truk Pengangkutan Sampah Di Kota Sumbawa Besar Shift li Menggunakan Gvrp," Jurnal UJMC (Unisda Journal of Mathematics and Computer Science), vol. 4, no. 2, pp. 15-23, 2018.
[13] L. O. M. Zulfiqar, R. R. Isnanto, and O. D. Nurhayati, "Optimal distribution route planning based on collaboration of dijkstra and sweep algorithm," Proceedings of 2018 10th International Conference on Information Technology and Electrical Engineering: Smart Technology for Better Society, ICITEE 2018, pp. 371-375, 2018, doi: 10.1109/ICITEED.2018.8534753.
[14] M. R. Wayahdi, S. H. N. Ginting, and D. Syahputra, "Greedy, A-Star, and Dijkstra's Algorithms in Finding Shortest Path," International Journal of Advances in Data and Information Systems, vol. 2, no. 1, pp. 45-52, Feb. 2021, doi: 10.25008/ijadis.v2i1.1206.
[15] L. Joni Erawati Dewi, "PENCARIAN RUTE TERPENDEK TEMPAT WISATA DI BALI DENGAN MENGGUNAKAN ALGORITMA DIJKSTRA," 2010.
[16] J. J. Siang, Riset Operasi Dalam Pendekatan Algoritmis. 2014.
[17] K. Hermanto and T. D. Ermayanti, "Analisa Optimasi Rute Transportasi Antar Jemput Siswa Menggunakan Model CGVRP dan Algoritma Dijkstra di SDIT Darus Sunnah," Unisda Journal of Mathematics and Computer Science (UJMC), vol. 5, no. 2, pp. 19-28, 2019, doi: 10.52166/ujmc.v5i2.1653.
[18] K. Hermanto, E. Ruskartina, and Gunawan, "Application of the CGVRP model and the 01 linear programming determines the RASKIN (subsidized rice) distribution route," Journal of Physics: Conference Series, vol. 1778, no. 1, pp. 0-8, 2021, doi: 10.1088/1742-6596/1778/1/012002.
[19] K. Hermanto, I. Adiasa, S. Altarisi, R. Rabani, and M. Amirul, "Rute Usulan Pendistribusian LPG Menggunakan Model Clustered Generalized Vehicle Routing Problem (CGVRP) dan Algoritma Dijkstra," Performa: Media Ilmiah Teknik Industri, vol. 19, no. 1, pp. 27-36, 2020, doi: 10.20961/performa.19.1.41858.

