Energy Efficient System for Wireless Sensor Networks using Modified RECHS Protocol

Mr. Nikhil Vilasrao Deshmukh Department of Electronics Engineering K.I.T.'s College of Engineering Kolhapur, India *n.deshmukh83@gmail.com*

Mr. Y. M. Patil Department of Electronics Engineering K.I.T.'s College of Engineering Kolhapur, India *ymp2002@rediffmail.com*

Abstract— The area of wireless sensor networks (WSNs) is one of the emerging and fast growing fields in the scientific world. This has brought about developing low cost, low-power and multi-function sensor nodes. Prolonged network lifetime, scalability, node mobility and load balancing are important requirements for many WSN applications. Clustering the sensor nodes is an effective technique to achieve these goals. Cluster-based routing protocol is currently a hot research in wireless sensor network. In this paper, we have added additional criteria for the selection of cluster heads in a Redundant and Energy-efficient Cluster head Selection Protocol (RECHS) and compared results with Energy Aware Low Energy Adaptive Clustering Hierarchy (EA-LEACH) protocol. This modified RECHS significantly increases the lifetime, reliability of the network. Simulation results show that comparison between two methods (Modified RECHS and EA- LEACH) for LEACH protocol on the basis of network lifetime (stability period), number of cluster heads are present per round, number of alive node are present per round and throughput of data transfer in the network.

Keywords- Wireless sensor network; Clustering algorithm; Redundant and energy-efficient clusterhead selection(RECHS); Stability period; LEACH protocol; EA-LEACH protocol

I. INTRODUCTION

A wireless sensor network is a collection of sensor nodes with limited power supply and constrained computational and transmission capability. Due to the limited transmission and computational ability, and high density of sensor nodes, forwarding of data packets takes place in multi-hop data transmission. Therefore routing in wireless sensor networks has been an important area of research in the past few years.

The sensor nodes run on non-rechargeable batteries, so along with efficient routing the network should be energy efficient with efficient utilization of the resources and hence this is an important research concern. Advances in wireless technologies and evolution of low cost sensor nodes have led to introduction of low power wireless sensor networks. Due to multiple functions and ease of deployment of the sensor nodes it can be used in various applications such as target tracking, environment monitoring, health care, forest fire detection, inventory control, energy management, surveillance and reconnaissance, and so on[1]. The main responsibility of the sensor nodes in a network is to forward the collected information from the source to the sink for further operations, but the resource limitations [2], unreliable links between the sensor nodes in combination with the various application demands of different applications make it a difficult task to design an efficient routing algorithm in wireless sensor networks.

Considering the unique features of low power wireless sensor networks, routing in WSN is much more challenging compared to traditional wireless networks such as ad-hoc networks [3, 4]. It has its own characteristics: such a nodes energy cannot be added, a large number of data redundancies in networks. AL-Karki et.al [3] classified the existing routing protocols in wireless sensor networks in two different perspectives, network structure and protocol operation. Hierarchical routing protocols were proposed to increase the scalability of the network and make the network energy efficient through node clustering. Clustering algorithms are recognized as the most effective organization. Many clustering algorithms in various contexts have been proposed [5,6,7,8].

In this paper, we have added additional criteria for the selection of cluster heads in a Redundant and Energy-efficient Clusterhead Selection Protocol (RECHS) and compared results with Energy Aware Low Energy Adaptive Clustering Hierarchy (EA-LEACH) protocol[9] that significantly increases the lifetime, reliability of the network.

In the Modified RECHS, selection criteria for cluster heads are weighted by remaining energy of a node and the average hops from all the ordinary nodes to the cluster heads in the network as well as we have added other criteria to select random number of cluster heads (within 1 to 6).

II. RELATED WORKS

We have studied a number of clustering algorithms in various contexts. Heinzelman et al have proposed a distributed algorithm for micro sensor networks in which the sensors select themselves as clusterheads with some probability and broadcast their decisions [5]. This algorithm allowed only 1hop clusters to be formed and had brought the convenience of communication. For the nodes which are far away the clusterhead need considerable energy to transmit data to clusterhead, such nodes will soon run out of battery power and premature be death. The lifetime of the network will reduce, and it is not conducive to expand the sensor networks.

Heinzelman et al have also proposed LEACH-C (LEACH-Centralized)[5]. It uses a centralized algorithm to form clusters. A non-autonomous cluster head selection is again the main disadvantage of this algorithm. Moreover, LEACH-C requires location information of all nodes of the network. Location information requires additional communication among the nodes. Lindsey et al present an interesting chainbased algorithm to solve the problem of collecting data from a micro sensor network [10]. M. J. Handy et al have proposed a deterministic clusterhead selection algorithm [11]. There is a modification of the threshold equation T (n). This algorithm only considered the selection of the majority of clusterheads, but for the problems of that some clusterheads are selected unreasonable and clusterheads are far away from the cluster nodes are not resolved. Then some Multi-hop routing clustering algorithms are proposed [6,12,13,14]. The way of cluster node transmit data to clusterhead is not single-hop, but multi-hop by other cluster nodes. Multi-hop clustering routing algorithm solves the problem from the cluster nodes which far away clusterhead to the clusterhead in Communication.

The algorithm proposed in this paper uses conditional probability and the number of hops to choose clusterhead and redundant clusterhead. In addition to this, we have also changed the criteria of selection of cluster heads at each round. The clusterheads selected are not random but are chosen depending upon the available energies of the nodes. In other words, out of the total number of nodes, those nodes which have maximum energies available have a higher priority to become the clusterheads. So, the clusterheads are chosen as evenly as possible which can only be achieved if at each round base station the available energy of each and every node. On the one hand, clusterheads are distributed in a more reasonable position, energy consumption will be decreased, and the network lifetime will be extended. On the other hand, the new protocol would solve the problem that the clusterhead fails or be damaged by attackers. The reliability and security of network have been improved.

III. MODIFIED REDUNDANT AND ENERGY EFFICIENT CLUSTERHEAD SELECTION PROTOCOL ALGORITHM (MODIFIED RECHS)

We will propose a Modified Redundant and Energy efficient Clusterhead Selection algorithm. In Modified RECHS, we first select initial clusterheads using random selection and optimal numbers (within 1-6) algorithms. After that we select initial redundant clusterheads and then compare the two nodes. The node which has better performance will be selected as clusterhead and other node will be as redundant clusterhead in current round.

A. Selection of Initial Clusterhead

In RECHS, we use a random selection algorithm with a condition that adds two parameters: Ecurrent and Emax. Ecurrent is node remaining energy, Emax is node initial energy. For a node S, S generates a random number rand (s) from 0-1, rand (s) multiply by a factor which representing the remaining energy level of a node be as the new random number rand '(s):

rand '(s) = rand(s)*(1- $E_{current} / E_{max})$

There is inverse relationship between remaining energy of a node and new random number. The probability of a node to become a clusterhead is greater if rand '(s) is less than the threshold T(s). If rand '(s) is bigger than threshold T(s), node S will be as the ordinary node.

B. Additional Criteria for Selection of Clusterheads in RECHS algorithm

Neither very less nor very large number of clusters are required to have a stable network. The number of clusters always has an optimal value at which the network is most stable. It is found from the study of [15] that the optimal value of clusters is given by

$$k_{opt} = \frac{\sqrt{N}}{\sqrt{2\Pi}} \sqrt{\frac{\mathcal{E}_{fs}}{\mathcal{E}_{mp}}} \frac{M}{d_{toBS}^2}$$

where,

k _{opt}	is optimal number of clusters(cluster heads),
Ń	is the number of nodes,
М	is the length of the region,
ϵf_s and ϵ_r	are the characteristics of the transmitter
d _{toBS}	is the distance between the base station and clusterhead

Along with this condition, we introduce a modification by restricting number of clusterheads to within 1 to 6 so that reduction in energy consumption of the nodes in the wireless sensor network is possible.

C. Selection of Initial redundant cluster head

Besides the initial clusterhead, we will select a second node in the cluster that the node has more energy than other nodes except the initial clusterhead in the cluster as the initial redundant clusterhead.Initial clusterhead and initial redundant clusterhead both have chance to be clusterhead.

D. Determination of cluster head and redundant clusterhead

In RECHS, we compare initial redundant clusterhead with initial clusterhead in their remaining energy and the sum of hops from other nodes to the two nodes.

- Formulae to calculate weights of energy and hops: $C_{k1} = red-hop/(red-hop + init-hop);$ $C_{k2} = init-hop/(red-hop + init-hop);$ $C_{e1} = E(CHvec(ii))/(E(CHvec(ii))+E(RCHvec(ii)));$ $C_{e2} = E(RCHvec(ii))/(E(CHvec(ii))+E(RCHvec(ii)));$
- Formulae to determine cluster head and redundant clusterhead: $C_1 = C_{k1} * C_{e1};$ $C_2 = C_{k2} * C_{e2};$

By comparison, we choose the node with optimal performance as the clusterhead in this round, and the rest node will become the redundant clusterhead.e.g., if $C_1 > C_2$ initial clusterhead is clusterhead (within cluster) for current round otherwise redundant clusterhead is clusterhead for current round.

IV. SIMULATION RESULTS

To validate the performance of modified RECHS, we will compare parameters of EA-LEACH and Modified RECHS as:

A. Performance Parameters of EA-LEACH and Modified RECHS:

MATLAB Simulator: 100-node random network; Eo=0.5J per node; Eelec = 50nJ/bit; camp = 100pJ/bit/m2; Nbits= 4000 bits; Base Station:(50,150); Area= 100 m2



Figure 1. Simulation result of the sample networks with base station

MATLAB Simulator: 100-node random network; Eo=0.5J per node; Eelec = 50nJ/bit; camp = 100pJ/bit/m2; Nbits= 4000 bits; Base Station:(50,150); Area= 100 m2



Figure 2. Simulation result of the sample networks with base station

B. Simulation results of Comparison of EA-LEACH and Modified RECHS:

Comparison on the basis of Number of Alive Node Vs Round:



Figure 3. Simulation result of stability period

In above figure result shows that in both methods, number of alive node is 100 till end of 670th round. After first node dead, other node dies very slowly as round increases in modified RECHS method where as in EA-LEACH node dies very fast.

Comparison on the basis of Residual Energy Vs Round:



Figure 4. Simulation result of residual energy of system with number of rounds

In above figure, result shows that In both method, Residual energy decays as round increases. In modified RECHS method Residual energy is greater than zero even after end of 800th round.

Comparison on the basis of Number of Clusterhead Vs Round:



Figure 5. Simulation result of number of clusterheads with number of rounds

In above figure, result shows that in RECHS method, Number of clusterhead per round mostly lies between 3 to 6. In EA-LEACH method, Number of clusterhead per round mostly lies between 2 to 6.

Comparison on the basis of Number of Packet Sent To BS Vs Round:



Figure 6. Simulation result of Packets received at the base station with number of rounds

In above figure, result shows that number of Packet sent to Base Station per round is nearly same in both method. Number

EA-LEACH EA-LEACH: First node died at 759 round EA-LEACH: Communication possible till at 777 round

Result Shows that

ompare to EA-LEACH method.

RECHS: First node died at 672 round RECHS: Communication possible till 821 round

Comparison on the basis of Network Lifetime:

V. CONCLUSION

of Packet sent to Base station is more in RECHS method as

Modified RECHS has better performance than EA-LEACH in terms of throughput and Network life time. Network stability period is more in case of EA-LEACH. In EA-LEACH method, we utilized full energy of all nodes for data transfer. This is not possible with RECHS method. As per application requirement, we can modify LEACH protocol by one of these methods to get optimum output from wireless sensor network.

REFERENCES

- [1] Jennifer Yick, Biswanath Mukherjee, and Dipak Ghosal. Wireless sensor network survey. Comput. Netw., 52(12):2292– 2330, August 2008.
- [2] Kamalrulnizam Abu Bakar Marjan Radi, Behnam Dezfouli and Malrey Lee. Multipath outing in wireless sensor networks: Survey and research challenges. MDPI Sensors, 12(1):650– 685, January 2012.
- [3] J. N. Al-karaki and A. E. Kamal. Routing techniques in wireless sensor networks: A survey. IEEE Wireless Communications, 11(6):6–28, December 2004.
- [4] Kemal Akkaya and Mohamed Younis. A survey on routing protocols for wireless sensor networks. Ad Hoc Networks, 3:325–349, 2005.
- [5] W. R. Heinzelman, A. P. Chandrakasan, and H. Balakrishnan, "Energy-efficient communication protocol for wireless microsensor networks", in Proceedings of the 33rd Hawaii International Conference on System Sciences (HICSS-33), January 2000.
- [6] S. Bandyopadhyay and E. J. Coyle, "An energy efficient hierarchical clustering algorithm for wireless sensor networks," in Proceedings of INFOCOM 2003, April 2003.
- [7] H. Su and X. Zhang, "Energy-efficient clustering system model and reconfiguration schemes for wireless sensor networks", in proc. of the 40th Conference on Information Sciences and Systems (CISS 2006), March 2006.
- [8] Hao Wang, Pu Tu, Ping Wang and Jian Yang, "A Redundant and Energy-Efficient Clusterhead Selection Protocol for Wireless Sensor Network", Conference on Communication Software and Networks, 2010.
- [9] Rahul Chhabra, Rahul Sharma, Mukul Bhardvaj, Pooja Garg "Energy Optimization In Routing Protocols Of Wireless Sensor Network", 2nd National Conference in Intelligent Computing & Communication Organized by Dept. of IT, GCET, Greater Noida, INDIA.
- [10] S. Lindsey, C. Raghavendra, K. Sivalingam, "Data gathering in sensor networks using the energy*delay metric", Proc. of the IPDPS Workshop on Issues in Wireless Networks and Mobile Computing, April 2001.
- [11] M.J.Handy, M.Haase, D.Timmermann, "Low energy adaptive clustering hierarchy with deterministic cluster-head selection", Mobile and Wireless Communications Network, 2002. 4th International Workshop on, 2002: 368- 372.
- [12] L. Qing, Q. Zhu, and M. Wang. Design of a distributed energyefficient clustering algorithm for heterogeneous wireless sensor networks. Computer Communications, 29(12):2230--2237, August 2006.

- [13] M.Younis, M.Youssef, K.Arisha "Energy-aware routing in cluster-based sensor networks", In Proceedings of MASCOTS 2002, 2002:129—136.
- [14] V.P.Mhatre, C. Rosenberg and D.Kofman and N.Shroff "A Minimum Cost Heterogeneous Sensor Networks with a lifetime Constraint". In Proceedings of IEEE Transactions on Mobile Computing, Vol.4, No.1. 2005.
- [15] Shio kumar Singh, M.P. Singh, and D.K. Singh, "A Survey of Energy-Efficient Hierarchical Cluster Based Routing in Wireless Sensor Networks", International Journal of Advanced Networking and Applications Volume:02, Issue:02,Pages:570-580,2010