Computer Science Faculty Publications and Presentations

College of Engineering and Computer Science

MMZDA: Enabling Social Welfare Maximization in Cross-Silo Federated Learning

Jianan Chen
Purdue University

Qin Hu
Purdue University

Honglu Jiang
The University of Texas Rio Grande Valley

Recommended Citation

Chen, Jianan, Qin Hu, and Honglu Jiang. "MMZDA: Enabling Social Welfare Maximization in Cross-Silo Federated Learning." arXiv preprint arXiv:2202.08362 (2022).

MMZDA: Enabling Social Welfare Maximization in Cross-Silo Federated Learning

Jianan Chen, Qin Hu, and Honglu Jiang

arXiv:2202.08362v2 [cs.GT] 2 Mar 2022

Abstract—As one of the typical settings of Federated Learning (FL), cross-silo FL allows organizations to jointly train an optimal Machine Learning (ML) model. In this case, some organizations may try to obtain the global model without contributing their local training, lowering the social welfare. In this paper, we model the interactions among organizations in cross-silo FL as a public goods game for the first time and theoretically prove that there exists a social dilemma where the maximum social welfare is not achieved in Nash equilibrium. To overcome this social dilemma, we employ the Multi-player Multi-action Zero-Determinant (MMZD) strategy to maximize the social welfare. With the help of the MMZD, an individual organization can unilaterally control the social welfare without extra cost. Since the MMZD strategy can be adopted by all organizations, we further study the scenario where multiple organizations jointly adopt the MMZD strategy and form an MMZD Alliance (MMZDA). We prove theoretically that the MMZDA strategy strengthens the control of the maximum social welfare. Experimental results validate that the MMZD strategy is effective in maximizing the social welfare and the MMZDA can achieve a larger maximum value.

Index Terms—Federated learning, public goods game, zero-determinant strategy, social welfare, game theory

I. INTRODUCTION

IN Federated Learning (FL), clients cooperatively train a Machine Learning (ML) model with their decentralized datasets under the coordination of a central server [1]. One of the typical settings of FL is cross-silo FL [2] where a neutral third-party agent acts as the central server and clients are a group of organizations, aiming to jointly train an optimal ML model for their respective use. In this case, these organizations are also the owners of the global model and can utilize the well-trained global model to further process tasks for their own interests.

An optimal global model with high performance requires the organizations in cross-silo FL to collaborate efficiently so as to bring considerable benefits to all participants, which can be regarded as the maximization of the social welfare. In fact, there are many studies on optimizing the social welfare in cross-silo FL by directly improving the model performance [3]–[6], increasing the convergence speed [7], reducing the communication cost [8], protecting privacy [9]–[11] and security [12], etc.

Jianan Chen and Qin Hu (Corresponding Author) are with the Department of Computer and Information Science, Indiana University - Purdue University Indianapolis, IN, USA. Email: jc144@iu.edu, qinhu@iu.edu

Honglu Jiang is with the Department of Informatics and Engineering Systems, The University of Texas Rio Grande Valley, Brownsville, TX, USA. Email: honglu.jiang@utrgv.edu

However, since every organization in cross-silo FL can obtain the final global model regardless of its contribution, the well-trained model becomes a public good, which is non-excludable and non-rivalrous for all organizations [13]. This leads to selfish behaviors that some organizations may only consider their own interests via inactively participating in local training to obtain the final global model for free or at a lower cost. The spread of this behavior can result in huge loss of the social welfare, and then none of the organizations can get the optimal model, which compromises the long-term stability and sustainability of cross-silo FL.

Most of the existing studies motivate organizations to fully contribute to cross-silo FL by designing an incentive mechanism [13]–[17]. However, this requires extra negotiation costs since organizations need to reach a consensus on the mechanism in advance, and demands additional running costs where a distributed algorithm runs over all organizations, clearly adding more burden to organizations. Recently, game-theoretical approach is applied to investigate participation behaviors in cross-silo FL [18], which needs additional information transmission between the server and organizations, and thus may cause potential privacy issues.

In this paper, we take a brand-new approach using the Multi-player Multi-action Zero-Determinant (MMZD) strategy [19] to maximize the social welfare in cross-silo FL without causing additional costs and transmission for organizations. Moreover, the MMZD Alliance (MMZDA) formed by multiple MMZD players is able to control the maximum social welfare more effectively. Another outstanding advantage of our methods is that they can control the social welfare and be applied to any cross-silo FL scenario no matter what strategies or actions other organizations perform. In summary, our contributions include (a preliminary version of this paper is presented in ICASSP 2022 [20]):

- We model the interactions among organizations in cross-silo FL as a public goods game for the first time, focusing on the organization's strategy rather than designing an extra mechanism to solve the social welfare maximization problem.
- We reveal the existence of the social dilemma in cross-silo FL by mathematical proof for the first time, which demonstrates the adverse effect of selfish behaviors in cross-silo FL in the view of game theory. This can be used as a theoretical basis for exploring organizations' behaviors in cross-silo FL.
- We overcome the social dilemma by employing the

MMZD strategy from the perspectives of individual organization and alliance. Specifically, any organization can unilaterally maximize the social welfare, which ensures the social welfare in cross-silo FL at a certain level and maintains the stability of the system.

- We further study the scenario in which multiple organizations adopt the same MMZD strategy, forming the MMZDA. We have mathematically proved that the maximum social welfare controlled by the MMZDA can get improved. This approach can also extend the application domains of the MMZD.
- Experiments prove the effectiveness of the MMZD strategy in maximizing the social welfare. The maximum value of social welfare is able to be enlarged by the MMZDA strategy.

The rest of the paper is organized as follows. The related work about cross-silo FL is summarized in Section II. In Section III, we formulate the cross-silo FL game to model the interactions among organizations in cross-silo FL, and further discover the social dilemma in the game. We propose a method based on MMZD strategy for individual organization to control the social welfare in Section IV. Section V studies the MMZD strategy employed by multiple organizations, namely MMZDA, and validates that the MMZD alliance can enlarge the maximum value of social welfare compared to the individual MMZD. Simulation results are reported in Section VI, followed by the conclusion in Section VII.

II. RELATED WORK

Existing works related to cross-silo FL can be classified into three categories: the optimization of aggregation algorithm, security and privacy protection, and the incentive mechanism design.

Aggregation algorithms are developed to enhance the performance of cross-silo FL. Based on the original FedAvg algorithm [1], various algorithms to improve the convergence speed, accuracy, and security were proposed in cross-silo FL setting. Marfoq et al. introduced practical algorithms to design an averaging policy under a decentralized model for achieving the fastest convergence [7]. And Huang et al. proposed FedAMP to overcome non-iid challenges [3]. Zhang et al. reduced the encryption and communication overhead caused by additively homomorphic encryption, which lowered the cost of aggregation as well [8].

Security and privacy protection issues also received attention in cross-silo FL. Heikkil et al. combined additively homomorphic secure summation protocols with differential privacy to guarantee strict privacy for individual data subjects in the cross-silo FL setting [9]. Li et al. proposed a brand-new one-shot algorithm that can flexibly achieve differential privacy guarantees [21]. Jiang et al. designed FLASHE, an optimized homomorphic encryption scheme to meet requirements of semantic security and additive homomorphism in cross-silo FL [12]. Chu et al. proposed a federated estimation method to accurately estimate the fairness of a model, namely avoiding the bias of data, without infringing the data privacy of any party [22].

Many studies used incentive mechanisms to encourage organizations to participate in training since the performance of cross-silo FL is affected by organizations' behaviors. Tang et al. formulated a social welfare maximization problem and proposed an incentive mechanism for cross-silo FL to address the organization heterogeneity and the public goods characteristics [13]. Many incentive mechanisms based on auction [14], contract [15], [16], and pricing [17] were proposed for FL. Unfortunately, they cannot be adapted to the cross-silo FL. First of all, most of the existing incentive mechanisms focus on encouraging more clients to participate in FL, instead of improving the performance of the global model from the perspective of social welfare. Secondly, the global model in cross-silo FL has the non-exclusive nature of public goods, leading to potential free-riding behaviors, which is not considered in the existing incentive mechanisms. Last but not least, current incentive mechanisms usually require participants or servers to spend additional computing resources. As a result, complex designs and additional computing costs can bring a burden on the organizations.

Recently, Zhang et al. explored the long-term participation of organizations in cross-silo FL by deploying a game theory approach [18], where a strategy and an algorithm were displayed for organizations to reduce free riders and increase the amount of local data for global training. However, the proposed algorithm requires extra information interactions between the organizations and server, which leads to the risk of information leakage.

In light of the above analysis, the following aspects distinguish our work from the existing approaches. First, our study reveals the social dilemma in the cross-silo FL. Second, we are committed to improving the final performance of the FL model, namely the social welfare. Moreover, we adopt the MMZD strategy without additional cost to control the maximum value of social welfare, and we use the MMZD alliance to further expand the ability to control the maximum value of social welfare.

III. SYSTEM MODEL

We consider a cross-silo FL scenario with a set of organizations, denoted as $\mathcal{N} = \{1, 2, ..., N\}$. All organizations rely on a central server to collaboratively conduct global model training for a specific task, where each of them has their own data for local training. The goal of organizations is to obtain an optimal global model, minimizing the loss based on all datasets. The central server collects the results of local model updates from all organizations, aggregates to obtain the global model, and then distributes it to everyone for the next round of training.

In each round of local training, every organization performs K iterations of model training. We denote the number of global communication rounds for aggregation as r. For the current task, the action of organization $i \in \mathcal{N}$, denoted as $y_i \in \{0, 1, ..., r\}$, represents the number of communication rounds it participates in the task. Then, $\mathbf{y} = (y_1, ..., y_i, ..., y_N)$ denotes the action vector of all organizations. Here we assume that all organizations in this cross-silo FL may participate

in fewer global aggregations due to laziness or selfishness, but they do not carry out malicious attacks, such as model poisoning attack.

According to the cross-silo FL model, all organizations get the same model in return. Inspired by [13], we define the revenue of organization i as:

$$\Phi_i(\mathbf{y}) = m_i(\chi_0 - \chi(\mathbf{y})), \tag{1}$$

where m_i (in dollars per unit of precision function) denotes the unit revenue of organization i by using the returned final model, χ_0 denotes the precision of the untrained model, and $\chi(\mathbf{y})$ denotes the precision of the trained global model after the actions of organizations in the action vector $\mathbf{y}$. Specifically, $\chi(\mathbf{y})$ can be modeled as

$$\chi(\mathbf{y}) = \frac{\theta_0}{\theta_1 + K \sum_{i \in N} y_i}, \tag{2}$$

with positive coefficients θ_0 and θ_1 [23] being derived based on the loss function, neural network, and local datasets. In particular, we have $\chi_0 = \frac{\theta_0}{\theta_1}$. The revenue of each organization is proportional to the difference between the expected loss after r rounds aggregation (i.e., $\chi(\mathbf{y})$) and the minimum expected loss (i.e., χ_0) [13]. As the number of total participation rounds increases, the marginal decrease of the difference reduces.

We define the cost of organization i as:

$$\Psi_i(y_i) = C_p^i + C_m^i. \tag{3}$$

The cost is composed of the organization's computation cost C_p^i and its communication cost C_m^i. On the one hand, the computation cost $C_p^i = \beta_i K y_i$, where β_i is a positive parameter, denoting the computation cost of each iteration in organization i's local training[1]. On the other hand, the communication cost C_m^i is defined as a parameter since we assume that every organization uploads its updates in each global aggregation round. If it chooses not to participate in global aggregation, it will submit a zero vector as updates.

Then the utility of organization i is defined as the difference between its revenue and cost:

$$U^i(\mathbf{y}) = \Phi_i(\mathbf{y}) - \Psi_i(y_i). \tag{4}$$

According to previous statements, we model the interactions among organizations as a *cross-silo FL game*.

Definition III.1. (Cross-silo FL game). In the cross-silo FL game, organizations participating in cross-silo FL act as players, where their actions and utilities are y_i and $U^i(\mathbf{y})$, respectively.

The cross-silo FL game can be iterative since these organizations in cross-silo FL usually cooperate for a long time to finish multiple FL tasks. Each game round in the cross-silo FL game corresponds to a certain FL task. Moreover, the social welfare in cross-silo FL game can be denoted as the total utility of all organizations i, namely $\Sigma_{i=1}^N U^i(\mathbf{y})$. In the

cross-silo FL game, we find that the social dilemma occurs if $\Phi_i(\mathbf{y}) - C_p^i < 0$, which can be summarized as below.

Theorem 1. (Social dilemma). If $\Phi_i(\mathbf{y}) - C_p^i < 0$, there exists a social dilemma in the cross-silo FL game.

Proof. The social dilemma forms when the Nash equilibrium is not the point of maximum social welfare. First, we study the Nash equilibrium of the cross-silo FL game. Referring to (4), we can derive the derivative of U^i as $m_i \frac{K\theta_0}{(\theta_1 + K\Sigma y_i)^2} - \beta_i K$. Given $\Phi_i(\mathbf{y}) - C_p^i < 0$, we have

$$m_i \frac{Ky_i\theta_0}{(\theta_1 + Ky_i)\theta_1} - \beta_i Ky_i < 0,$$

which leads to

$$m_i \frac{K\theta_0}{(\theta_1 + K\Sigma y_i)^2} < m_i \frac{K\theta_0}{(\theta_1 + Ky_i)\theta_1} < \beta_i K.$$

Thus, the derivative of U^i is negative, and the utility function decreases monotonically with y_i. The Nash equilibrium strategy of each organization is $y_i = 0$, so the Nash equilibrium point is $\mathbf{y}^{NE} = (0, 0, ..., 0)$. Noted that, there is a natural and necessary premise of FL that the utility of the well-trained model must be positive. Thus, we prove that the point $\mathbf{y}^r = (y_i = r, i \in \mathcal{N})$ with the social welfare

$$\sum_{i=1}^N U^i(\mathbf{y}^r) = \sum \Psi_i(\mathbf{y}^r) - \sum C_p^i - \sum C_m^i > 0,$$

higher than that in the Nash equilibrium point

$$\sum_{i=1}^N U^i(\mathbf{y}^{NE}) = -\sum C_m^i < 0.$$

So the social dilemma exists if $\Phi_i(\mathbf{y}) - C_p^i < 0$. $\square$

The condition $\Phi_i(\mathbf{y}) - C_p^i < 0$ in the above theorem indicates that if any organization $i \in \mathcal{N}$ only trains the local model using its local dataset, the utility is negative. In fact, this condition strengthens organizations' motivation to participate in global training in cross-silo FL game.

Thus, there is a social dilemma in the cross-silo FL game and problem follows. If the organization only pursues its own interests and does not participate in the communication round, it will lead to a decrease in the accuracy of the global model, which in turn leads to low social welfare. So it is of great importance to optimize the public goods in the cross-silo FL game, i.e., the accuracy of the global model, which is actually equivalent to the social welfare.

For reference, we summarize key notations used in the system model in Table I.

IV. SOCIAL WELFARE MAXIMIZATION BY MMZD

According to the analysis above, we can see that the underlying cause of the social dilemma is selfishness, leading to the loss of all organizations, namely the low social welfare. Aiming to solve this problem, we resort to the Multi-player Multi-action Zero-Determinant (MMZD) strategy for the social welfare maximization in this section. In each game round, any organization can choose the action

---

[1]As [24] shows, $\beta_i = \frac{\alpha_i}{2} f_i^2 D_i S_i$, where $\frac{\alpha_i}{2}$ is the effective capacitance coefficient of organization i's computing chipset, f_i denotes the calculation processing capacity, D_i denotes the number of data units, and S_i denotes the number of CPU cycles required by organization i to process one data unit.

TABLE I
KEY NOTATIONS.

Notation	Meaning
N	The number of organizations
K	The number of local training iterations for every organization
r	The number of global communication rounds for aggregation
y_i	The number of communication rounds organization i participates in the current task
$\mathbf{y}$	The action vector of all organizations
Φ_i	The revenue of organization i
m_i	The unit revenue of organization i by using the final model
$\chi(\mathbf{y})$	The precision of the trained global model with the corresponding action vector $\mathbf{y}$
Ψ_i	The cost of organization i
C_p^i	The computation cost of organization i
C_m^i	The communication cost of organization i
β_i	The computation cost of each iteration in organization i's local training
$U^i(\mathbf{y})$	The utility of organization i with the corresponding action vector $\mathbf{y}$

$y_i \in \{0, 1, ..., r\}$, so there are $(r + 1)^N$ possible outcomes for each game round. We assume that the organizations have one-round memory since a long-memory player has no priority against others with short memory [19]. Fig. 1 describes an example of the cross-silo FL game with two organizations and three actions, i.e., $N = 2$ and $r = 2$, in which all possible outcomes can be denoted as $(y_1, y_2) \in \{(0,0), (0,1), (0,2), (1,0), (1,1), (1,2), (2,0), (2,1), (2,2)\}$. For arbitrary organization $i \in \mathcal{N}$, its mixed strategy $\mathbf{p}^i$ is defined as:

$$\mathbf{p}^i = [p_{1,0}^i, p_{1,1}^i, ..., p_{1,r}^i, p_{2,0}^i, ..., p_{j,g}^i, ..., p_{(r+1)^N,r}^i]^T, \quad (5)$$

where $p_{j,g}^i (j \in \{1, 2, ..., (r+1)^N\}, g \in \{0, 1, ..., r\})$ represents the probability of organization i choosing action $y_i = g$ in the current game round and other organizations choosing the same actions as j-th outcome of the previous game round. As presented in Fig. 1, for the previous outcome $(y_1, y_2) = (0, 2)$, the conditional probability of organization 1 to select action $y_1 = 1$ is $p_{3,1}^1$ and the conditional probability of organization 2 to select action $y_2 = 0$ is $p_{3,0}^2$. In addition, the corresponding utility vector $\mathbf{u}^i$ is denoted as:

$$\mathbf{u}^i = [u_{1,0}^i, u_{1,1}^i, ..., u_{1,r}^i, u_{2,0}^i, ..., u_{j,g}^i, ..., u_{(r+1)^N,r}^i]^T, \quad (6)$$

where each utility $u_{j,g}^i$ of organization i choosing action $y_i = g$ in the j-th outcome can be calculated by $u_{j,g}^i = U^i(\mathbf{y}^{(j,g)})$, with $\mathbf{y}^{(j,g)}$ denoting the action vector $\mathbf{y}$ corresponding to the j-th outcome but $y_i = g$. Fig. 2(a) presents the strategy vectors and utility vectors based on the example shown in Fig. 1.

In the cross-silo FL model, an organization's current move depends only on its last action and the action vector $\mathbf{y}$ in the last game round. We can construct a Markov matrix $\mathbf{M} = [M_{vw}]_{(r+1)^N \times (r+1)^N}$, with each element M_{vw} denoting the one-step transition probability from state v to w. Fig. 2(b) shows the Markov matrix $\mathbf{M}$ for the case of $r = 2$ and $N = 2$. For example, the element $p_{3,1}^1 p_{3,0}^2$ is at the 3rd row and 4th column in Fig. 2, which represents the possibility of transitioning from the 3rd-outcome (0,2) in the previous round to the 4th-outcome (1,0) in the current round. Then, we define $\mathbf{M}' \equiv \mathbf{M} - \mathbf{I}$ where $\mathbf{I}$ is an identity matrix. And we assume
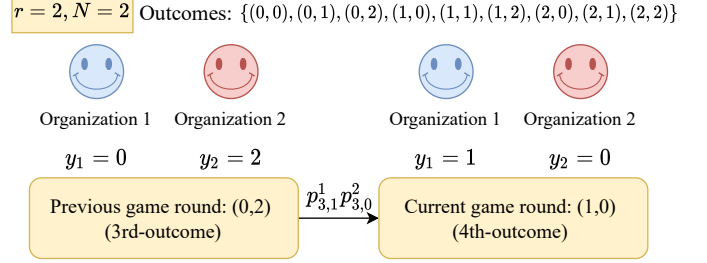

Fig. 1. Illustration of the cross-silo FL game with two organizations and three actions ($N = 2$ and $r = 2$).

the stationary vector of $\mathbf{M}$ is π. Given $\pi^T \mathbf{M} = \pi^T$, we can draw that $\pi^T \mathbf{M}' = 0$. According to Cramer's rule, we have

$$Adj(\mathbf{M}')\mathbf{M}' = \det(\mathbf{M}')\mathbf{I} = 0, \quad (7)$$

where $Adj(\mathbf{M}')$ denotes the adjugate matrix of $\mathbf{M}'$. Thus, it can be noted that every row of $Adj(\mathbf{M}')$ is proportional to π. Hence for any vector $\mathbf{a} = (a_1, a_2, ..., a_{(r+1)^N})^T$, we can draw $\pi^T \cdot \mathbf{a} = \det(\mathbf{p}^1, ..., \mathbf{p}^N, \mathbf{a})$ according to [25].

In particular, when $\mathbf{a} = \mathbf{u}^i$, then organization i's expected utility in the stationary state is:

$$E^i = \frac{\pi^T \cdot \mathbf{u}^i}{\pi^T \cdot \mathbf{1}} = \frac{\det(\mathbf{p}^1, ..., \mathbf{p}^N, \mathbf{u}^i)}{\det(\mathbf{p}^1, ..., \mathbf{p}^N, \mathbf{1})}, \quad (8)$$

which makes a linear combination of all organizations' expected utilities yielding the following equation:

$$\sum_{x=1}^{N} \alpha_x E^x + \alpha_0 = \frac{\det(\mathbf{p}^1, ..., \mathbf{p}^N, \sum_{x=1}^{N} \alpha_x \mathbf{u}^x + \alpha_0 \mathbf{1})}{\det(\mathbf{p}^1, ..., \mathbf{p}^N, \mathbf{1})}. \quad (9)$$

In the above equation, α_0 and $\alpha_x, x \in \mathcal{N}$, are constants for the linear combination. Moreover, after some elementary column operations on $\det(\mathbf{p}^1, \mathbf{p}^2, ..., \mathbf{p}^N, \mathbf{a})$, in which there can be a certain column controlled by a certain organization. Fig. 2(c) displays the determinant of $\pi^T \cdot \mathbf{a}$ based on the example in Fig. 1, in which the first column is solely determined by organization 1. Thus, when organization i chooses a strategy that satisfies

$$\hat{\mathbf{p}}^i = \phi(\sum_{x=1}^{N} \alpha_x \mathbf{u}^x + \alpha_0 \mathbf{1}), \quad (10)$$

where ϕ is a non-zero constant and $\hat{\mathbf{p}}^i$ is under the control of organization i, the corresponding column of $\hat{\mathbf{p}}^i$ and the last column of $\det(\mathbf{p}^1, ..., \mathbf{p}^N, \sum_{x=1}^{N} \alpha_x \mathbf{u}^x + \alpha_0 \mathbf{1})$ will be proportional. And (9) can be converted to:

$$\sum_{x=1}^{N} \alpha_x E^x + \alpha_0 = 0. \quad (11)$$

We further study the social welfare maximization problem with the help of the MMZD in this circumstance. Take organization 1 performing the MMZD strategy as an example. According to (11), by setting $\alpha_x = 1, x \in \mathcal{N}$, the social welfare can be calculated as $\sum_{x=1}^{N} E^x = -\alpha_0$. Thus, the issue of maximizing

(a)

$$\mathbf{p}^1 = [p^1_{1,0}, p^1_{1,1}, p^1_{1,2}, p^1_{2,0}, p^1_{2,1}, p^1_{2,2}, p^1_{3,0}, p^1_{3,1}, p^1_{3,2}, p^1_{4,0}, p^1_{4,1}, p^1_{4,2}, p^1_{5,0}, p^1_{5,1}, p^1_{5,2}, p^1_{6,0}, p^1_{6,1}, p^1_{6,2}, p^1_{7,0}, p^1_{7,1}, p^1_{7,2}, p^1_{8,0}, p^1_{8,1}, p^1_{8,2}, p^1_{9,0}, p^1_{9,1}, p^1_{9,2}]$$

$$\mathbf{p}^2 = [p^2_{1,0}, p^2_{1,1}, p^2_{1,2}, p^2_{2,0}, p^2_{2,1}, p^2_{2,2}, p^2_{3,0}, p^2_{3,1}, p^2_{3,2}, p^2_{4,0}, p^2_{4,1}, p^2_{4,2}, p^2_{5,0}, p^2_{5,1}, p^2_{5,2}, p^2_{6,0}, p^2_{6,1}, p^2_{6,2}, p^2_{7,0}, p^2_{7,1}, p^2_{7,2}, p^2_{8,0}, p^2_{8,1}, p^2_{8,2}, p^2_{9,0}, p^2_{9,1}, p^2_{9,2}]$$

$$\mathbf{u}^1 = [u^1_{1,0}, u^1_{1,1}, u^1_{1,2}, u^1_{2,0}, u^1_{2,1}, u^1_{2,2}, u^1_{3,0}, u^1_{3,1}, u^1_{3,2}, u^1_{4,0}, u^1_{4,1}, u^1_{4,2}, u^1_{5,0}, u^1_{5,1}, u^1_{5,2}, u^1_{6,0}, u^1_{6,1}, u^1_{6,2}, u^1_{7,0}, u^1_{7,1}, u^1_{7,2}, u^1_{8,0}, u^1_{8,1}, u^1_{8,2}, u^1_{9,0}, u^1_{9,1}, u^1_{9,2}]$$

$$\mathbf{u}^2 = [u^2_{1,0}, u^2_{1,1}, u^2_{1,2}, u^2_{2,0}, u^2_{2,1}, u^2_{2,2}, u^2_{3,0}, u^2_{3,1}, u^2_{3,2}, u^2_{4,0}, u^2_{4,1}, u^2_{4,2}, u^2_{5,0}, u^2_{5,1}, u^2_{5,2}, u^2_{6,0}, u^2_{6,1}, u^2_{6,2}, u^2_{7,0}, u^2_{7,1}, u^2_{7,2}, u^2_{8,0}, u^2_{8,1}, u^2_{8,2}, u^2_{9,0}, u^2_{9,1}, u^2_{9,2}]$$

(b)

$$\mathbf{M} = \begin{bmatrix}
p^1_{1,0}p^2_{1,0} & p^1_{1,0}p^2_{1,1} & p^1_{1,0}p^2_{1,2} & p^1_{1,1}p^2_{1,0} & p^1_{1,1}p^2_{1,1} & p^1_{1,1}p^2_{1,2} & p^1_{1,2}p^2_{1,0} & p^1_{1,2}p^2_{1,1} & p^1_{1,2}p^2_{1,2} \\
p^1_{2,0}p^2_{2,0} & p^1_{2,0}p^2_{2,1} & p^1_{2,0}p^2_{2,2} & p^1_{2,1}p^2_{2,0} & p^1_{2,1}p^2_{2,1} & p^1_{2,1}p^2_{2,2} & p^1_{2,2}p^2_{2,0} & p^1_{2,2}p^2_{2,1} & p^1_{2,2}p^2_{2,2} \\
p^1_{3,0}p^2_{3,0} & p^1_{3,0}p^2_{3,1} & p^1_{3,0}p^2_{3,2} & p^1_{3,1}p^2_{3,0} & p^1_{3,1}p^2_{3,1} & p^1_{3,1}p^2_{3,2} & p^1_{3,2}p^2_{3,0} & p^1_{3,2}p^2_{3,1} & p^1_{3,2}p^2_{3,2} \\
p^1_{4,0}p^2_{4,0} & p^1_{4,0}p^2_{4,1} & p^1_{4,0}p^2_{4,2} & p^1_{4,1}p^2_{4,0} & p^1_{4,1}p^2_{4,1} & p^1_{4,1}p^2_{4,2} & p^1_{4,2}p^2_{4,0} & p^1_{4,2}p^2_{4,1} & p^1_{4,2}p^2_{4,2} \\
p^1_{5,0}p^2_{5,0} & p^1_{5,0}p^2_{5,1} & p^1_{5,0}p^2_{5,2} & p^1_{5,1}p^2_{5,0} & p^1_{5,1}p^2_{5,1} & p^1_{5,1}p^2_{5,2} & p^1_{5,2}p^2_{5,0} & p^1_{5,2}p^2_{5,1} & p^1_{5,2}p^2_{5,2} \\
p^1_{6,0}p^2_{6,0} & p^1_{6,0}p^2_{6,1} & p^1_{6,0}p^2_{6,2} & p^1_{6,1}p^2_{6,0} & p^1_{6,1}p^2_{6,1} & p^1_{6,1}p^2_{6,2} & p^1_{6,2}p^2_{6,0} & p^1_{6,2}p^2_{6,1} & p^1_{6,2}p^2_{6,2} \\
p^1_{7,0}p^2_{7,0} & p^1_{7,0}p^2_{7,1} & p^1_{7,0}p^2_{7,2} & p^1_{7,1}p^2_{7,0} & p^1_{7,1}p^2_{7,1} & p^1_{7,1}p^2_{7,2} & p^1_{7,2}p^2_{7,0} & p^1_{7,2}p^2_{7,1} & p^1_{7,2}p^2_{7,2} \\
p^1_{8,0}p^2_{8,0} & p^1_{8,0}p^2_{8,1} & p^1_{8,0}p^2_{8,2} & p^1_{8,1}p^2_{8,0} & p^1_{8,1}p^2_{8,1} & p^1_{8,1}p^2_{8,2} & p^1_{8,2}p^2_{8,0} & p^1_{8,2}p^2_{8,1} & p^1_{8,2}p^2_{8,2} \\
p^1_{9,0}p^2_{9,0} & p^1_{9,0}p^2_{9,1} & p^1_{9,0}p^2_{9,2} & p^1_{9,1}p^2_{9,0} & p^1_{9,1}p^2_{9,1} & p^1_{9,1}p^2_{9,2} & p^1_{9,2}p^2_{9,0} & p^1_{9,2}p^2_{9,1} & p^1_{9,2}p^2_{9,2}
\end{bmatrix}$$

(c)

$$\pi^T \mathbf{a} = \det(\mathbf{p}^1, \mathbf{p}^2, \mathbf{a})$$

$$= \det \begin{bmatrix}
-1+p^1_{1,0} & p^1_{1,0}p^2_{1,1} & p^1_{1,0}p^2_{1,2} & p^1_{1,1}p^2_{1,0} & p^1_{1,1}p^2_{1,1} & p^1_{1,1}p^2_{1,2} & p^1_{1,2}p^2_{1,0} & p^1_{1,2}p^2_{1,1} & a_1 \\
-1+p^1_{2,0} & -1+p^1_{2,0}p^2_{2,1} & p^1_{2,0}p^2_{2,2} & p^1_{2,1}p^2_{2,0} & p^1_{2,1}p^2_{2,1} & p^1_{2,1}p^2_{2,2} & p^1_{2,2}p^2_{2,0} & p^1_{2,2}p^2_{2,1} & a_2 \\
-1+p^1_{3,0} & p^1_{3,0}p^2_{3,1} & -1+p^1_{3,0}p^2_{3,2} & p^1_{3,1}p^2_{3,0} & p^1_{3,1}p^2_{3,1} & p^1_{3,1}p^2_{3,2} & p^1_{3,2}p^2_{3,0} & p^1_{3,2}p^2_{3,1} & a_3 \\
p^1_{4,0} & p^1_{4,0}p^2_{4,1} & p^1_{4,0}p^2_{4,2} & -1+p^1_{4,1}p^2_{4,0} & p^1_{4,1}p^2_{4,1} & p^1_{4,1}p^2_{4,2} & p^1_{4,2}p^2_{4,0} & p^1_{4,2}p^2_{4,1} & a_4 \\
p^1_{5,0} & p^1_{5,0}p^2_{5,1} & p^1_{5,0}p^2_{5,2} & p^1_{5,1}p^2_{5,0} & -1+p^1_{5,1}p^2_{5,1} & p^1_{5,1}p^2_{5,2} & p^1_{5,2}p^2_{5,0} & p^1_{5,2}p^2_{5,1} & a_5 \\
p^1_{6,0} & p^1_{6,0}p^2_{6,1} & p^1_{6,0}p^2_{6,2} & p^1_{6,1}p^2_{6,0} & p^1_{6,1}p^2_{6,1} & -1+p^1_{6,1}p^2_{6,2} & p^1_{6,2}p^2_{6,0} & p^1_{6,2}p^2_{6,1} & a_6 \\
p^1_{7,0} & p^1_{7,0}p^2_{7,1} & p^1_{7,0}p^2_{7,2} & p^1_{7,1}p^2_{7,0} & p^1_{7,1}p^2_{7,1} & p^1_{7,1}p^2_{7,2} & -1+p^1_{7,2}p^2_{7,0} & p^1_{7,2}p^2_{7,1} & a_7 \\
p^1_{8,0} & p^1_{8,0}p^2_{8,1} & p^1_{8,0}p^2_{8,2} & p^1_{8,1}p^2_{8,0} & p^1_{8,1}p^2_{8,1} & p^1_{8,1}p^2_{8,2} & p^1_{8,2}p^2_{8,0} & -1+p^1_{8,2}p^2_{8,1} & a_8 \\
p^1_{9,0} & p^1_{9,0}p^2_{9,1} & p^1_{9,0}p^2_{9,2} & p^1_{9,1}p^2_{9,0} & p^1_{9,1}p^2_{9,1} & p^1_{9,1}p^2_{9,2} & p^1_{9,2}p^2_{9,0} & p^1_{9,2}p^2_{9,1} & a_9
\end{bmatrix}$$

Fig. 2. The strategy vectors, utility vectors, Markov matrix, and determinant of $\pi^T \cdot \mathbf{a}$ after elementary transformations in the cross-silo FL game example with two organizations and three actions shown in Fig. 1. (a) The strategy vectors and the corresponding utility vectors of the cross-silo FL game. (b) The Markov matrix $\mathbf{M}$ of the cross-silo FL game. (c) After several elementary column operations on $\det(\mathbf{p}^1, \mathbf{p}^2, \mathbf{a})$, the dot product of the stationary vector π and an arbitrary vector $\mathbf{a} = (a_1, a_2, \ldots, a_9)^T$ is equal to $\det(\mathbf{p}^1, \mathbf{p}^2, \mathbf{a})$, where the first columns $\hat{\mathbf{p}}^1$ is only controlled by organization 1.

the social welfare is equivalent to the following optimization problem:

$$\min \alpha_0,$$
$$s.t. \begin{cases} 0 \le p^1_{j,g} \le 1, j \in \{1, 2, \ldots, (r+1)^N\}, g \in \{0, \ldots, r\}, \\ \hat{\mathbf{p}}^1 = \phi(\sum_{x=1}^N \mathbf{u}^x + \alpha_0 \mathbf{1}), \\ \phi \ne 0. \end{cases}$$

We denote $u^x_k, k \in \{1, 2, \ldots, (r+1)^{N+1}\}$ as the kth element in $\mathbf{u}^x$, then we can solve the above optimization problem by considering the following two cases:

A. $\phi > 0$

To meet the constraint $p^1_{j,g} \ge 0$, we can get the lower bound of α_0 as follows:

$$\alpha_{0min} = \max(\Lambda_k), \forall k \in \{1, 2, \ldots, (r+1)^{N+1}\},$$
$$\Lambda_k = \begin{cases} -\sum_{x=1}^N u^x_k - \frac{1}{\phi}, k = 1, 2, \ldots(r+1)^{N-1}, \\ -\sum_{x=1}^N u^x_k, k = (r+1)^{N-1} + 1, \ldots, (r+1)^{N+1}. \end{cases}$$

To meet the constraint $p^1_{j,g} \le 1$, we can get the upper bound of α_0 as follows:

$$\alpha_{0max} = \min(\Lambda_l), \forall l \in \{(r+1)^{N+1} + 1, \ldots, 2(r+1)^{N+1}\},$$
$$\Lambda_l = \Lambda_{k+(r+1)^N}$$
$$= \begin{cases} -\sum_{x=1}^N u^x_k, & k = 1, 2, \ldots(r+1)^{N-1}, \\ -\sum_{x=1}^N u^x_k + \frac{1}{\phi}, & k = (r+1)^{N-1} + 1, \ldots, (r+1)^{N+1}. \end{cases}$$

Only if $\alpha_{0min} \le \alpha_{0max}$, can α_0 have a feasible solution, which is equivalent to $\max(\Lambda_k) \le \min(\Lambda_l), \forall k \in \{1, 2, \ldots, (r+1)^{N+1}\}, \forall l \in \{(r+1)^N + 1, \ldots, 2(r+1)^{N+1}\}$. If there exists $\phi > 0$ satisfying the above constraint, we can obtain the minimum value of α_0 as follow:

$$\alpha_{0min} = \max\{-\sum_{x=1}^N u^x_1 - \frac{1}{\phi}, \ldots, -\sum_{x=1}^N u^x_{(r+1)^{N-1}} - \frac{1}{\phi},$$
$$-\sum_{x=1}^N u^x_{(r+1)^{N-1}+1}, \ldots, -\sum_{x=1}^N u^x_{(r+1)^{N+1}}\}. \quad (12)$$

B. $\phi < 0$

Similarly, when $p_{j,g}^1 \geq 0$, we have $\alpha_{0min} = \max(\Lambda_l), \forall l \in \{(r+1)^{N+1} + 1, ..., 2(r+1)^{N+1}\}$; considering $p_{j,g}^1 \leq 1$, we have $\alpha_{0max} = \min(\Lambda_k), \forall k \in \{1, 2, ..., (r+1)^{N+1}\}$. In addition, α_0 is feasible only when $\alpha_{0min} \leq \alpha_{0max}$, i.e., $\max(\Lambda_l) \leq \min(\Lambda_k), \forall k \in \{1, 2, ..., (r+1)^{N+1}\}, \forall l \in \{(r+1)^{N+1} + 1, ..., 2(r+1)^{N+1}\}$. Finally, we can get the following result:

$$
\alpha_{0min} = \max\{-\sum_{x=1}^N u_1^x, ..., -\sum_{x=1}^N u_{(r+1)^{N-1}}^x,
$$
$$
-\sum_{x=1}^N u_{(r+1)^{N-1}+1}^x + \frac{1}{\phi}, ..., -\sum_{x=1}^N u_{(r+1)^{N+1}}^x + \frac{1}{\phi}\}. \quad (13)
$$

In summary, by (12) and (13), organization 1 can unilaterally set the expected social welfare $\sum_{x=1}^N E^x$ with the MMZD strategy $\mathbf{p}^1$ meeting $\hat{\mathbf{p}}^1 = \phi(\sum_{x=1}^N \mathbf{u}^x + \alpha_0 \mathbf{1})$, with each element of $\mathbf{p}^1$ calculated by:

$$
p_h^1 = \begin{cases} \sum_{x=1}^N u_h^x + \alpha_{0min} + 1, \\ \qquad\qquad h = 1, 2, ..., (r+1)^{N-1}, \\ \sum_{x=1}^N u_h^x + \alpha_{0min}, \\ \qquad\qquad h = (r+1)^{N-1} + 1, ..., (r+1)^{N+1}, \end{cases}
$$

where p_h^1 denotes the h-th element in $\mathbf{p}^1$.

V. SOCIAL WELFARE MAXIMIZATION BY MMZD ALLIANCE

In the previous section, we proved that a single organization is able to maximize the social welfare within a range. According to (10), every organization can deploy the MMZD strategy to control the social welfare. Therefore, it's possible that multiple organizations utilize the MMZD strategy at the same time. So in this section, we explore multiple organizations that play the MMZD strategy to form an alliance in the cross-silo FL game. We call them MMZD Alliance (MMZDA) organizations (denoted as $\mathcal{A}$), assuming that all MMZDA organizations use the same MMZD strategy to prevent the low social welfare. Besides, we define the other organizations as outsider organizations (denoted as $\mathcal{N} \backslash \mathcal{A}$). Outsider organizations may inactively participate in communication rounds and just want to get model trained by other organizations for free.

The current challenge is how MMZDA controls the maximum value of social welfare. Specifically, will the MMZDA have stronger control and enlarge the maximum social welfare compared with the individual MMZD strategy?

For the sake of convenience, we assume that organization $i \in \mathcal{A}, |\mathcal{A}| = N^{\mathcal{A}}$ is an alliance, which performs the same MMZD strategy as the leader organization $a \in \mathcal{A}$. As for an outsider organization $j \in \mathcal{N} \backslash \mathcal{A}, |\mathcal{N} \backslash \mathcal{A}| = N - N^{\mathcal{A}}$, it may not participate in any communication round.

In this new scenario based on the MMZDA, we pay more attention to strategies and behaviors rather than the organizations themselves. Since the MMZDA members take the same actions, we treat them as an entity represented by the leader organization a. In each game round, any outsider organization or alliance leader (organization a) can choose the action

$y_i \in \{0, 1, ..., r\}$, so there are $(r+1)^{N-N^{\mathcal{A}}}$ possible outcomes for each game round. We assume that the organizations have one-round-memory. And we define $\mathcal{N}^c = \mathcal{N} \backslash \mathcal{A} \cup \{a\}$ as the players of the cross-silo FL game based on MMZDA. For arbitrary organization $i \in \mathcal{N}^c$, its mixed strategy $\mathbf{q}^i$ is defined as:

$$
\mathbf{q}^i = [q_{1,0}^i, q_{1,1}^i, ..., q_{1,r}^i, q_{2,0}^i, ..., q_{j,g}^i, ..., q_{(r+1)^{N-N^{\mathcal{A}}+1},r}^i]^T, \quad (14)
$$

where $q_{j,g}^i (j \in \{1, 2, ..., (r+1)^{N-N^{\mathcal{A}}+1}\}, g \in \{0, 1, ..., r\})$ represents the probability of organization i choosing action $y_i = g$ in the current game round conditioning on the j-th outcome of the previous game round. In addition, the corresponding utility vector $\mathbf{v}^i$ is denoted as:

$$
\mathbf{v}^i = [v_{1,0}^i, v_{1,1}^i, ..., v_{1,r}^i, v_{2,0}^i, ..., v_{j,g}^i, ..., v_{(r+1)^{N-N^{\mathcal{A}}+1},r}^i]^T, \quad (15)
$$

where each utility $v_{j,g}^i$ of organization i choosing action $y_i = g$ in the j-th outcome can be calculated by $v_{j,g}^i = U^i(\mathbf{y}^{(j,g)})$, with $\mathbf{y}^{(j,g)}$ denoting the action vector $\mathbf{y}$ corresponding to the j-th outcome but $y_i = g$. Moreover, if $i = a$, then $v_{j,g}^a = \sum_{x \in \mathcal{A}} U^x(\mathbf{y}^{(j,g)})$. In Section IV, we perform that the linear combination of all organizations' expected utilities can be represented as (8) and (9). Similarly, we can draw that

$$
\sum_{x \in \mathcal{N}^c} \gamma_x F^x + \gamma_0
$$
$$
= \frac{\det(\mathbf{q}^1, ..., \mathbf{q}^i, ..., \mathbf{q}^{N-N^{\mathcal{A}}+1}, \sum_{x \in \mathcal{N}^c} \gamma_x \mathbf{v}^x + \gamma_0 \mathbf{1})}{\det(\mathbf{q}^1, ..., \mathbf{q}^i, ..., \mathbf{q}^{N-N^{\mathcal{A}}+1}, \mathbf{1})}, \quad (16)
$$

where $\gamma_0 \in \mathbb{R}$ and $\gamma_x \in \mathbb{R}, x \in \mathcal{N}^c$, are constants. Thus, when organization i chooses a strategy that satisfies $\hat{\mathbf{q}}^i = \phi(\sum_{x \in \mathcal{N}^c} \gamma_x \mathbf{v}^x + \gamma_0 \mathbf{1})$, where ϕ is a non-zero constant and $\hat{\mathbf{q}}^i$ is under the control of organization i, the column related to $\hat{\mathbf{q}}^i$ and the last column of $\det(\mathbf{q}^1, ..., \mathbf{q}^i, ..., \mathbf{q}^N, \mathbf{v}^i)$ will be proportional. Then (9) can be converted to:

$$
\sum_{x \in \mathcal{N}^c} \gamma_x F^x + \gamma_0 = 0. \quad (17)
$$

In order to investigate the social welfare maximization problem by MMZDA, we rewrite the optimization problem as:

$$
\min \gamma_0,
$$
$$
s.t. \begin{cases} 0 \leq q_{j,g}^a \leq 1, \\ j \in \{1, 2, ..., (r+1)^{N-N^{\mathcal{A}}+1}\}, g \in \{0, ..., r\}, \\ \hat{\mathbf{q}}^a = \phi(\sum_{x \in \mathcal{N}^c} \mathbf{v}^x + \gamma_0 \mathbf{1}), \\ \phi \neq 0. \end{cases}
$$

Similar to Section IV, we denote $v_k^x, k \in \{1, 2, ..., (r+1)^{N-N^{\mathcal{A}}+2}\}$ as the kth element in $\mathbf{v}^x$, then we can solve this optimization problem by discussing these two situations:
(1) $\phi > 0$:
When $q_{j,g}^a \geq 0$, we have $\gamma_{0min} = \max(\Lambda_l), \forall l \in \{(r+1)^{N-N^{\mathcal{A}}+2} + 1, ..., 2(r+1)^{N-N^{\mathcal{A}}+2}\}$; given $q_{j,g}^a \leq 1$, we have $\gamma_{0max} = \min(\Lambda_k), \forall k \in \{1, 2, ..., (r+1)^{N-N^{\mathcal{A}}+2}\}$. In addition, γ_0 is feasible only when $\gamma_{0min} \leq \gamma_{0max}$, i.e., $\max(\Lambda_l) \leq \min(\Lambda_k), \forall k \in \{1, 2, ..., (r+1)^{N-N^{\mathcal{A}}+2}\}, \forall l \in$

$\{(r+1)^{N-N^{\mathcal{A}}+2}+1,...,2(r+1)^{N-N^{\mathcal{A}}+2}\}$. Finally, we can get the following result:

$$\gamma_{0min} = \max\{-\sum_{x\in\mathcal{N}^c}v_1^x - \frac{1}{\phi},...,-\sum_{x\in\mathcal{N}^c}v_{(r+1)^{N-N^{\mathcal{A}}}}^x - \frac{1}{\phi},$$
$$-\sum_{x\in\mathcal{N}^c}v_{(r+1)^{N-N^{\mathcal{A}}+1}}^x,...,-\sum_{x\in\mathcal{N}^c}v_{(r+1)^{N-N^{\mathcal{A}}+2}}^x\}. \quad (18)$$

(2) $\phi < 0$:
Given $q_{j,g}^a \geq 0$, we have $\gamma_{0min} = \max(\Lambda_l), \forall l \in \{(r+1)^{N-N^{\mathcal{A}}+2}+1,...,2(r+1)^{N-N^{\mathcal{A}}+2}\}$; while $q_{j,g}^a \leq 1$, we can get $\gamma_{0max} = \min(\Lambda_k), \forall k \in \{1,2,...,(r+1)^{N-N^{\mathcal{A}}+2}\}$. In addition, γ_0 is feasible only when $\gamma_{0min} \leq \gamma_{0max}$, i.e., $\max(\Lambda_l) \leq \min(\Lambda_k), \forall k \in \{1,2,...,(r+1)^{N-N^{\mathcal{A}}+2}\}, \forall l \in \{(r+1)^{N-N^{\mathcal{A}}+2}+1,...,2(r+1)^{N-N^{\mathcal{A}}+2}\}$. Finally, we can get the following result:

$$\gamma_{0min} = \max\{-\sum_{x\in\mathcal{N}^c}v_1^x,...,-\sum_{x\in\mathcal{N}^c}v_{(r+1)^{N-N^{\mathcal{A}}}}^x,$$
$$-\sum_{x\in\mathcal{N}^c}v_{(r+1)^{N-N^{\mathcal{A}}+1}}^x + \frac{1}{\phi},...,-\sum_{x\in\mathcal{N}^c}v_{(r+1)^{N-N^{\mathcal{A}}+2}}^x + \frac{1}{\phi}\}. \quad (19)$$

In summary, by (18) and (19), the alliance can unilaterally set the expected social welfare $\sum_{x\in\mathcal{N}^c}E^x$ with the MMZD strategy $\mathbf{q}^a$ meeting $\hat{\mathbf{q}}^a = \phi(\sum_{x\in\mathcal{N}^c}\mathbf{v}^x + \gamma_0\mathbf{1})$, with each element of $\mathbf{q}^a$ calculated by:

$$q_h^a = \begin{cases} \sum_{x=1}^N v_h^x + \gamma_{0min} + 1, \\ \quad h = 1,2,...,(r+1)^{N-N^{\mathcal{A}}}, \\ \sum_{x=1}^N v_h^x + \gamma_{0min}, \\ \quad h = (r+1)^{N-N^{\mathcal{A}}} + 1,...,(r+1)^{N-N^{\mathcal{A}}+2}, \end{cases}$$

where q_h^a denotes the h-th element in $\mathbf{q}^a$.

Theorem 2. The social welfare can achieve larger maximum value by MMZDA than that by single MMZD organization.

Proof. In the same cross-silo FL game, we can achieve a maximum social welfare $\sum_{x=1}^N E^x = -\alpha_0$ using MMZD strategy by an individual organization. Meanwhile, we are able to draw another maximum social welfare $\sum_{x\in\mathcal{N}^c}\gamma_x F^x = -\gamma_0$ by MMZDA.

A. $\phi > 0$

In this case, we have:

$$\alpha_{0min} = \max\{-\sum_{x=1}^N u_1^x - \frac{1}{\phi},...,-\sum_{x=1}^N u_{(r+1)^{N-1}}^x - \frac{1}{\phi},$$
$$-\sum_{x=1}^N u_{(r+1)^{N-1}+1}^x,...,-\sum_{x=1}^N u_{(r+1)^{N+1}}^x\}.$$
$$\gamma_{0min} = \max\{-\sum_{x\in\mathcal{N}^c}v_1^x - \frac{1}{\phi},...,-\sum_{x\in\mathcal{N}^c}v_{(r+1)^{N-N^{\mathcal{A}}}}^x - \frac{1}{\phi},$$
$$-\sum_{x\in\mathcal{N}^c}v_{(r+1)^{N-N^{\mathcal{A}}+1}}^x,...,-\sum_{x\in\mathcal{N}^c}v_{(r+1)^{N-N^{\mathcal{A}}+2}}^x\}.$$

The value of α_{0min} is the maximum value in the certain $(r+1)^{N+1}$ values, and we denote this candidate values

as set X_1. While the value of γ_{0min} generating from the $(r+1)^{N-N^{\mathcal{A}}+2}$ values, we denote this candidate values as set X_2. Note that the $(r+1)^{N+1}$ values cover all possible outcomes, but the $(r+1)^{N-N^{\mathcal{A}}+2}$ values do not cover all outcomes, since the MMZDA organizations employ the same strategy. Given a certain element like $-\sum_{x\in\mathcal{N}^c}v_k^x - \frac{1}{\phi} \in X_2, k \in \{1,2,...,N-N^{\mathcal{A}}\}$, as $v_{j,g}^a = \sum_{x\in\mathcal{A}}U^x(\mathbf{y}^{(j,g)})$, we have:

$$-\sum_{x\in\mathcal{N}^c}v_k^x - \frac{1}{\phi} = \sum_{x\in\mathcal{A}}U^x(\mathbf{y}^{(j,g)}) + -\sum_{x\in\mathcal{N}^c}v_k^x - \frac{1}{\phi}$$
$$= \sum_{x=1}^N U^x(\mathbf{y}^{(j\times(r+1)^{N^{\mathcal{A}-1}}-1,g)}) - \frac{1}{\phi} = -\sum_{x=1}^N u_{k'}^x - \frac{1}{\phi}, \quad (20)$$

where $k' = (r+1)^{N^{\mathcal{A}+k-1}}$. Otherwise, given an element like $-\sum_{x\in\mathcal{N}^c}v_k^x \in X_2, k \in \{N-N^{\mathcal{A}}+1,...,N-N^{\mathcal{A}}+2\}$, we can draw

$$-\sum_{x\in\mathcal{N}^c}v_k^x = \sum_{x\in\mathcal{A}}U^x(\mathbf{y}^{(j,g)}) + -\sum_{x\in\mathcal{N}^c}v_k^x$$
$$= \sum_{x=1}^N U^x(\mathbf{y}^{(j-(r+1)^{N-N^{\mathcal{A}-1}})\times(r+1)^{N^{\mathcal{A}-1}}+(r+1)^{N-1},g)})$$
$$= -\sum_{x=1}^N u_{k'}^x, (k' = (r+1)^{N^{\mathcal{A}+(k-N+N^{\mathcal{A}})-1}} + (r+1)^{N-1})$$

As deduced above, the elements in X_2 are all in X_1, so X_2 is a subset of X_1. Thus, $\alpha_{0min} \geq \gamma_{0min}$ holds true.

B. $\phi < 0$

When $\phi < 0$, the situation is quite similar.

$$\alpha_{0min} = \max\{-\sum_{x=1}^N u_1^x,...,-\sum_{x=1}^N u_{(r+1)^{N-1}}^x,$$
$$-\sum_{x=1}^N u_{(r+1)^{N-1}+1}^x + \frac{1}{\phi},...,-\sum_{x=1}^N u_{(r+1)^{N+1}}^x + \frac{1}{\phi}\}.$$
$$\gamma_{0min} = \max\{-\sum_{x\in\mathcal{N}^c}v_1^x,...,-\sum_{x\in\mathcal{N}^c}v_{(r+1)^{N-N^{\mathcal{A}}}}^x,$$
$$-\sum_{x\in\mathcal{N}^c}v_{(r+1)^{N-N^{\mathcal{A}}+1}}^x + \frac{1}{\phi},...,-\sum_{x\in\mathcal{N}^c}v_{(r+1)^{N-N^{\mathcal{A}}+2}}^x + \frac{1}{\phi}\}.$$

We denote the $(r+1)^{N+1}$ candidate values of α_{0min} as X_3. While the value of γ_{0min} generating from the $(r+1)^{N-N^{\mathcal{A}}+2}$ values, we denote this candidate values as X_4. We can prove that each value of X_4 can be found in X_3 as above. So $\alpha_{0min} \geq \gamma_{0min}$ holds true as well.

In summary, we can draw a conclusion that $\alpha_{0min} \geq \gamma_{0min}$, which is equivalent to $-\alpha_{0max} \leq -\gamma_{0max}$. So the social welfare can achieve larger maximization by MMZDA than single MMZD organization. $\square$

The above theorem and previous derivation prove that MMZDA further enhances the ability of non-selfish organizations to maximize the social welfare. If more organizations join the MMZDA, the upper boundary of controllable social welfare continues to increase, leading the entire system to be more stable.

VI. EXPERIMENTAL EVALUATION

In this section, we present the experimental results of our study for the MMZD individual (MMZD) strategy and MMZD Alliance (MMZDA) stategy in social welfare maximization. Generally, all experiments are implemented using Matlab R2021a on a laptop with 2.3 GHz Intel Core i5-8300H processor. In all experiments except for otherwise specification, we set $\phi = 0.01$, $K = 200$, $r = 33$. Parameters $\theta_0 = 23271.584$ and $\theta_1 = 50193.243$ are derived based on the simulation dataset [23]. For every control group with different strategy settings, we repeat the above experiment 100 times, and take the average value as the final expected social welfare.

A. Evaluation of the MMZD Strategy

First, we evaluate the performance of the MMZD strategy used by individual organization to maximize the social welfare based on simulation experiments. We set $N = 10$ since the number of organizations in cross-silo FL is usually small.
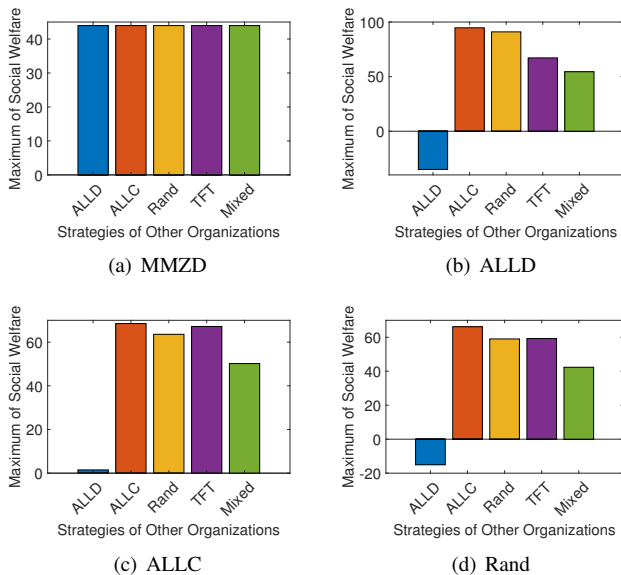

Fig. 3. The maximum of social welfare under different strategy combinations of organization 1 and others.

In order to verify the effectiveness of the MMZD strategy on maximizing the social welfare, we compare it with other five classical strategies, by simulating the entire cross-silo FL process for 20 rounds of game. In Fig. 3, organization 1 adopts MMZD, all-defection (ALLD) [26], all-cooperation (ALLC) [26], and random (Rand) [26] strategies. Other organizations use ALLD, ALLC, Rand, Tit-For-Tat (TFT) [27], and mixed (Mixed) strategies. Specifically, ALLD strategy is defined as: the organization does not perform local training at all, only submits a zero vector in the global aggregation. While ALLC strategy means that the organization participates in all r global aggregation with their local updates in every game round. Organizations which adopt Rand strategy randomly participate in global aggregation from 0 to r rounds with the probability of $\frac{1}{r+1}$. TFT strategy is defined as the organizations randomly choose the number of participating global aggregation from
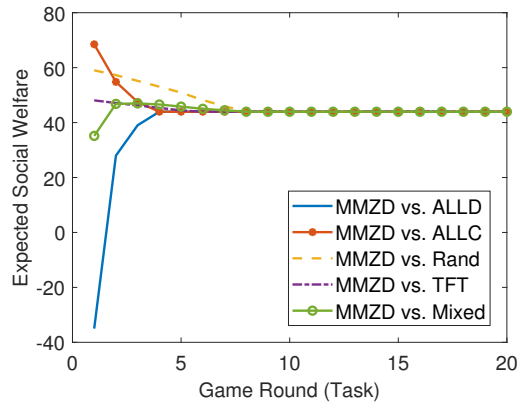

Fig. 4. Evolution of the expected social welfare.

0 to $\frac{\lfloor r \rfloor}{2}$ when the sum of global aggregation rounds in last game round is less than $\frac{Nr}{2}$, otherwise they randomly choose the number of participating global aggregation from $\frac{\lfloor r+1 \rfloor}{2}$ to r. We define the mixed strategy as adopting a specific strategy chosen from ALLC, ALLD, Rand, and TFT. By comparing Fig. 3(a) with the other three figures, we can find that the MMZD strategy can effectively control the social welfare. This can prove that free-riding behavior is avoided in some ways.

Fig. 4 plots the expected social welfare changes in each game round, as the organization 1 adopts the MMZD strategy and other organizations employ different strategies. Fig. 3(a) displays the final result in Fig. 4, which indicates that no matter what kind of strategies other organizations adopt, the social welfare finally converges to a fixed value, verifying the power of the proposed social welfare maximization game.

B. Evaluation of the MMZDA Strategy

In this subsection, we present the performance of the MMZDA to maximize the social welfare through a series of simulation experiments. And we compare it with the result of a single organization's MMZD strategy. Meanwhile, we also consider the relative maximum of social welfare in order to further analyze the control ability of MMZDA.

In Fig. 5(a), we set $N^{\mathcal{A}} = 4$ and randomly choose four organizations to form a MMZDA, with other parameters unchanged in order to compare with the previous experiment (Fig. 3). From this figure, it's clear that no matter what strategies other organizations adopt, the MMZDA strategy expands the maximum value of the social welfare, comparing with the MMZD strategy performed by a single organization. This experimental result verifies Theorem 2. That is, the social welfare can achieve larger maximum value by MMZDA than that by the single MMZD organization.

Based on the same initial settings, Figs. 5(b)-(f) display the evolution process of of the expected social welfare, which is two-fold. The red line represents that the MMZDA strategy is used, while the blue line shows the impact of single organization using the MMZD strategy on the maximum of social welfare. By comparison, we can find that as the number of game rounds increases, no matter what strategies other organizations adopt, the MMZDA strategy always enables the

maximum social welfare gradually converge to a larger fixed value, but it does not have a faster convergence speed.

(a) The maximum of social welfare under MMZDA and MMZD

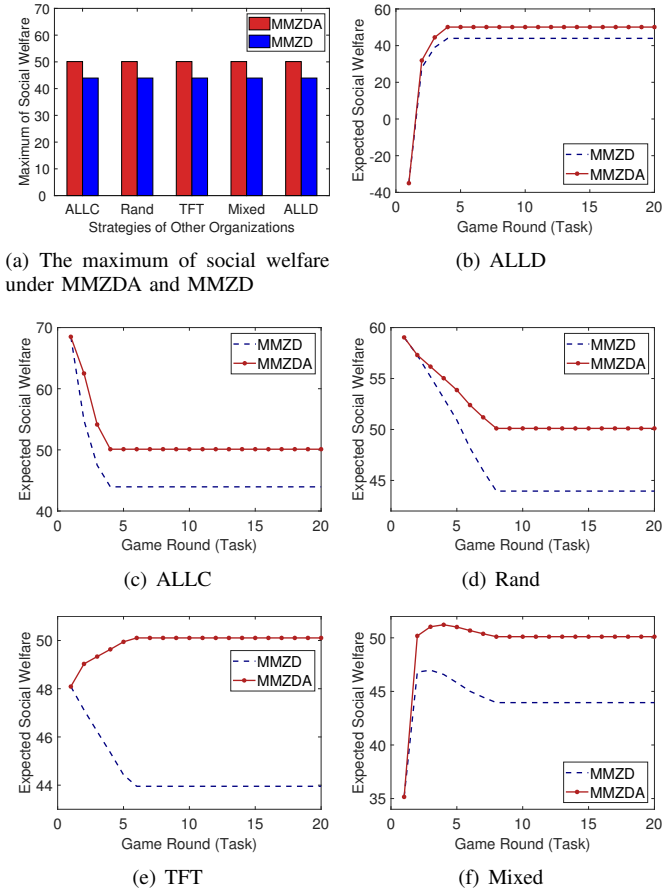(b) ALLD

(c) ALLC

(d) Rand

(e) TFT

(f) Mixed

Fig. 5. Evolution of the expected social welfare under different strategy combinations, under comparison between the MMZDA strategy and the MMZD individual strategy.

In Fig. 6, we set $N = 10$. Then in Fig. 6(a), we explore the changes of the expected social welfare as the number of organizations in the alliance increases when the total number of organizations remains unchanged. Whenever $N^{\mathcal{A}}$ takes a different value, we randomly generate the MMZDA organizations from this 10 organizations. In the histogram, we can conclude that when the total number of organizations does not change, the more organizations that join the MMZD alliance, the higher the maximum social welfare that can be controlled. This also confirms our analysis of the MMZDA strategy, because the increase of $N^{\mathcal{A}}$ expands the range of candidate values in (18) and (19), thereby increasing the maximum value of social welfare. Besides, in Fig. 6(b), we take the social welfare of all organizations participating in all communication rounds as the absolute maximum value of social welfare, and study the ratio of social welfare controlled by MMZDA to the former value. We call this ratio relative maximum of social welfare. In this case, $N^{\mathcal{A}}$ and the relative maximum of social welfare are also positively correlated. Together with Fig. 6(a), they show that when N is constant, MMZDA's ability to control the maximum value of social welfare increases as the number of MMZDA organizations increases.
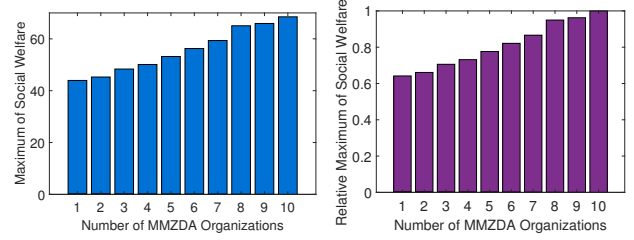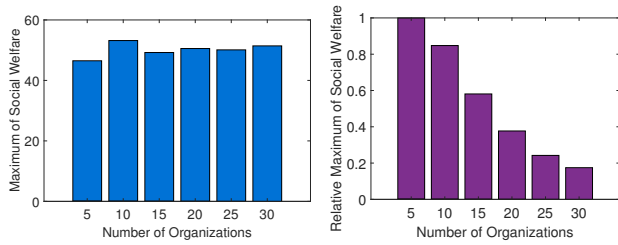

(a) The absolute maximum of social welfare

(b) The relative maximum of social welfare

Fig. 6. The absolute maximum of social welfare and the relative maximum of social welfare as the number of organizations in the MMZDA increases with the total number of organizations unchanged.

Correspondingly, we continue to investigate the impact of N while $N^{\mathcal{A}}$ does not change. In fact, the change of the total number of organizations N brings a series of differences, including adding new organizations' parameters, changing the organizations' local datasets, and then changing the coefficients θ_0 and θ_1, which also changes the utility function. For verification, we generate a new simulation dataset, and randomly selected N organizations in different situations. We continue to randomly select $N^{\mathcal{A}} = 5$ MMZDA organizations from N organizations. According to the new dataset, the corresponding coefficients θ_0 and θ_1 are generated to calculate the final social welfare value. For each distinct N, we repeat the process of randomly selecting N organizations 10 times. For each group of selected organizations, we repeat the process of randomly selecting five MMZDA organizations 10 times. Finally, we take the average value as the expected social welfare. As shown in Fig. 7(a), we found that simply changing N does not intuitively change the maximum value of social welfare, because the impact of newly joined organizations on social welfare is mutative. More specifically, when N is small or even close to $N^{\mathcal{A}}$ (i.e., the scenario of $N = 5$), the maximum value of social welfare is limited by the total number of organizations. When N is large (i.e., the scenario of $N = 25$), the small MMZD alliance reduces the ability to control the maximum of social welfare, and cannot bring a large increase of the maximum value. But in Fig. 7(b), N and the relative maximum of social welfare are negatively correlated. It reflects that the increase in N weakens MMZDA organizations' control over the maximum value of social welfare, although this does not mean a decrease in the absolute maximum value of social welfare.

In Fig. 8, we use the same dataset as the previous experiment (Fig. 7). In this experiment, we keep the ratio of the number of alliance organizations to the total number of organizations unchanged. Specifically, we set the number of alliance organizations to be $\frac{1}{3}$ of the total number of organizations. In Fig. 8(a), clearly, the maximum value of social welfare increases as N increases. This is because more organizations participate in cross-silo FL game, and the social welfare that can be increased when the proportion of MMZDA organizations remains unchanged. While Fig. 8(b) implies that under the same ratio, the relative maximum of social welfare fluctuates around 0.6 within a certain range. There is not

(a) The absolute maximum of social welfare

(b) The relative maximum of social welfare

Fig. 7. The absolute maximum of social welfare and the relative maximum of social welfare as the total number of organizations increases with the number of organizations in the MMZDA unchanged.
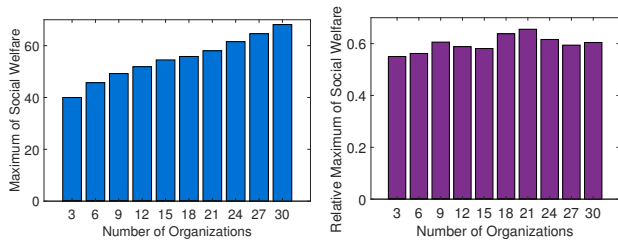
(a) The absolute maximum of social welfare

(b) The relative maximum of social welfare

Fig. 8. The absolute maximum of social welfare and the relative maximum of social welfare as the total number of organizations increases with the rate of MMZDA organizations unchanged.

much change overall, and the fluctuations come from the heterogeneity of the organizations. In fact, the control ability of MMZDA also depends on the utility vectors $\mathbf{v}$ of the alliance organizations. During the experiment, we randomly select the alliance organization to more objectively reflect the expected control ability of MMZDA.

VII. CONCLUSION

In this paper, we model the cross-silo FL game among organizations as a public goods game, revealing the social dilemma in cross-silo FL game theoretically. In order to overcome the social dilemma, we propose a brand-new method using the MMZD to solve the social welfare maximization problem. By the means of the MMZD, an individual organization can unilaterally control social welfare at a certain level, regardless of other organizations' strategies. Meanwhile, we explore the MMZDA consisting of multiple MMZD organizations, which further improves the control of the maximum social welfare. Moreover, our approaches can maintain the stability and sustainability of the system without extra cost. Simulation results prove that the MMZD strategy can efficiently and effectively control social welfare, which reduces the loss from selfish behaviors.

REFERENCES

[1] B. McMahan, E. Moore, D. Ramage, S. Hampson, and B. A. y Arcas, "Communication-efficient learning of deep networks from decentralized data," in *Artificial intelligence and statistics*. PMLR, 2017, pp. 1273–1282.

[2] P. Kairouz, H. B. McMahan, B. Avent, A. Bellet, M. Bennis, A. N. Bhagoji, K. Bonawitz, Z. Charles, G. Cormode, R. Cummings *et al.*, "Advances and open problems in federated learning," *arXiv preprint arXiv:1912.04977*, 2019.

[3] Y. Huang, L. Chu, Z. Zhou, L. Wang, J. Liu, J. Pei, and Y. Zhang, "Personalized cross-silo federated learning on non-iid data," in *Proceedings of the AAAI Conference on Artificial Intelligence*, vol. 35, no. 9, 2021, pp. 7865–7873.

[4] Y. Wang, Y. Tong, D. Shi, and K. Xu, "An efficient approach for cross-silo federated learning to rank," in *2021 IEEE 37th International Conference on Data Engineering (ICDE)*. IEEE, 2021, pp. 1128–1139.

[5] K. Nandury, A. Mohan, and F. Weber, "Cross-silo federated training in the cloud with diversity scaling and semi-supervised learning," in *ICASSP 2021-2021 IEEE International Conference on Acoustics, Speech and Signal Processing (ICASSP)*. IEEE, 2021, pp. 3085–3089.

[6] U. Majeed, L. U. Khan, and C. S. Hong, "Cross-silo horizontal federated learning for flow-based time-related-features oriented traffic classification," in *2020 21st Asia-Pacific Network Operations and Management Symposium (APNOMS)*. IEEE, 2020, pp. 389–392.

[7] O. Marfoq, C. Xu, G. Neglia, and R. Vidal, "Throughput-optimal topology design for cross-silo federated learning," *arXiv preprint arXiv:2010.12229*, 2020.

[8] C. Zhang, S. Li, J. Xia, W. Wang, F. Yan, and Y. Liu, "Batchcrypt: Efficient homomorphic encryption for cross-silo federated learning," in *2020 {USENIX} Annual Technical Conference ({USENIX}{ATC} 20)*, 2020, pp. 493–506.

[9] M. A. Heikkilä, A. Koskela, K. Shimizu, S. Kaski, and A. Honkela, "Differentially private cross-silo federated learning," *arXiv preprint arXiv:2007.05553*, 2020.

[10] Q. Li, B. He, and D. Song, "Practical one-shot federated learning for cross-silo setting," in *Proceedings of the Thirtieth International Joint Conference on Artificial Intelligence (IJCAI-21)*, 2021.

[11] G. Long, T. Shen, Y. Tan, L. Gerrard, A. Clarke, and J. Jiang, "Federated learning for privacy-preserving open innovation future on digital health," *arXiv preprint arXiv:2108.10761*, 2021.

[12] Z. Jiang, W. Wang, and Y. Liu, "Flashe: Additively symmetric homomorphic encryption for cross-silo federated learning," *arXiv preprint arXiv:2109.00675*, 2021.

[13] M. Tang and V. W. Wong, "An incentive mechanism for cross-silo federated learning: A public goods perspective," in *IEEE INFOCOM 2021-IEEE Conference on Computer Communications*. IEEE, 2021, pp. 1–10.

[14] M. Tang, H. Pang, S. Wang, L. Gao, J. Huang, and L. Sun, "Multidimensional auction mechanisms for crowdsourced mobile video streaming," *IEEE/ACM Transactions on Networking*, vol. 26, no. 5, pp. 2062–2075, 2018.

[15] N. Zhao, Y.-C. Liang, and Y. Pei, "Dynamic contract incentive mechanism for cooperative wireless networks," *IEEE Transactions on Vehicular Technology*, vol. 67, no. 11, pp. 10 970–10 982, 2018.

[16] S. Feng, D. Niyato, P. Wang, D. I. Kim, and Y.-C. Liang, "Joint service pricing and cooperative relay communication for federated learning," in *2019 International Conference on Internet of Things (iThings) and IEEE Green Computing and Communications (GreenCom) and IEEE Cyber, Physical and Social Computing (CPSCom) and IEEE Smart Data (SmartData)*. IEEE, 2019, pp. 815–820.

[17] Q. Shao, M. H. Cheung, and J. Huang, "Multimedia crowdsourcing with bounded rationality: A cognitive hierarchy perspective," *IEEE Journal on Selected Areas in Communications*, vol. 37, no. 7, pp. 1478–1488, 2019.

[18] N. Zhang, Q. Ma, and X. Chen, "Enabling long-term cooperation in cross-silo federated learning: A repeated game perspective," *IEEE Transactions on Mobile Computing*, 2022.

[19] X. He, H. Dai, P. Ning, and R. Dutta, "Zero-determinant strategies for multi-player multi-action iterated games," *IEEE Signal Processing Letters*, vol. 23, no. 3, pp. 311–315, 2016.

[20] J. Chen, Q. Hu, and H. Jiang, "Social welfare maximization in cross-silo federated learning," in *ICASSP 2022-2022 IEEE International Conference on Acoustics, Speech and Signal Processing (ICASSP)*. IEEE, 2022.

[21] Q. Li, B. He, and D. Song, "Practical one-shot federated learning for cross-silo setting," *arXiv preprint arXiv:2010.01017*, 2020.

[22] L. Chu, L. Wang, Y. Dong, J. Pei, Z. Zhou, and Y. Zhang, "Fedfair: Training fair models in cross-silo federated learning," *arXiv preprint arXiv:2109.05662*, 2021.

[23] X. Li, K. Huang, W. Yang, S. Wang, and Z. Zhang, "On the convergence of fedavg on non-iid data," *arXiv preprint arXiv:1907.02189*, 2019.

[24] N. H. Tran, W. Bao, A. Zomaya, M. N. Nguyen, and C. S. Hong, "Federated learning over wireless networks: Optimization model design and analysis," in *IEEE INFOCOM 2019-IEEE Conference on Computer Communications*. IEEE, 2019, pp. 1387–1395.

[25] W. H. Press and F. J. Dyson, "Iterated prisoner's dilemma contains strategies that dominate any evolutionary opponent," *Proceedings of the National Academy of Sciences*, vol. 109, no. 26, pp. 10409–10413, 2012.

[26] Q. Hu, S. Wang, P. Ma, X. Cheng, W. Lv, and R. Bie, "Quality control in crowdsourcing using sequential zero-determinant strategies," *IEEE Transactions on Knowledge and Data Engineering*, vol. 32, no. 5, pp. 998–1009, 2019.

[27] M. Nowak and K. Sigmund, "A strategy of win-stay, lose-shift that outperforms tit-for-tat in the prisoner's dilemma game," *Nature*, vol. 364, no. 6432, pp. 56–58, 1993.