ORIGINAL PAPER

Persistence and stability of a class of kinetic compartmental models

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Abstract

In this paper we show that the dynamics of a class of kinetic compartmental models with bounded capacities, monotone reaction rates and a strongly connected interconnection structure is persistent. The result is based on the chemical reaction network (CRN) and the corresponding Petri net representation of the system. For the persistence analysis, it is shown that all siphons in the Petri net of the studied model class can be characterized efficiently. Additionally, the existence and stability of equilibria are also analyzed building on the persistence and the theory of general compartmental systems. The obtained results can be applied in the analysis of general kinetic models based on the simple exclusion principle.

1 Introduction

Nonnegative systems form an important subclass within dynamical systems having the property that the nonnegative orthant is invariant with respect to the dynamics. The practical motivation for developing the theory of nonnegative systems is the fact that there are several application fields such as chemistry, biology, population and disease dynamics, where in many cases the state variables of the models in the original physical coordinates are nonnegative [18]. Compartmental models are used to describe the change of distribution of objects (e.g., molecules or particles) among different storage compartments in time [5]. Compartments can be physically distinct

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subsystems such as interconnected containers, but they can also represent disjoint states like different stages of diseases in the case of epidemic models [10]. Since the natural state variables in compartmental systems correspond to amounts of materials, numbers of molecules (or to their ratios, concentrations), these models belong to the nonnegative system class. The fundamental properties of compartmental models have been intensively studied in the literature. The observability, controllability, realizability and identifiability of compartmental systems are summarized in [11] focusing mainly on linear models. The analytic solution of linear compartmental ODEs is studied in [17] in a kinetic context. A fundamental reference on the qualitative analysis of a wide class of general nonlinear compartmental models is [22], where important results can be found on the structure of equilibria and stability.

It is known that most compartmental models can be represented in the form of kinetic systems also called chemical reaction networks (CRNs), where the dynamics can be formally realized by a set of reactions with appropriate complexes and reaction rates [15, 21]. Although kinetic models are originated from physical chemistry, they have been highly generalized in a mathematical sense (see, e.g. [3, 20, 29]), widening their application possibilities even to non-chemical processes as general descriptors of nonlinear dynamics. Chemical reaction network theory (CRNT) is a dynamically improving research field with strong results on the relations between the reaction graph structure and the qualitative properties of the kinetic dynamics [16]. Persistence analysis is a problem of central importance in CRNT, for instance, it is a key property for proving global asymptotic stability of complex balanced networks [2, 13, 14].

The application of discrete structures and graph theory is an essential tool in the modeling and analysis of chemical phenomena [4, 9, 12]. If the molecules and reactions in a kinetic system are tracked individually, the CRN can be considered as a discrete event system which is equivalent to a Petri net [6, 30]. A fundamental and theoretically deep result of CRNT is published in [7], where conditions are given for the persistence of continuous time CRNs using the graph structure of their Petri net representation and conserved quantities in the dynamics. These results were further generalized in [8] to time-dependent open systems, where reaction rates can be time-dependent and there are in and outflows.

The model class we study is related to the so-called simple exclusion principle known from the theory of Markov Processes [24] in the sense that particles can move along a directed graph (called the compartmental graph). It is also assumed that the capacity of the compartments is bounded. Therefore, transition is only possible if there are available particles in the donor compartment, and also free space in the recipient compartment. A well-known application of the simple exclusion principle is the class of ribosome flow models (RFMs) [33] capturing key features of the translation process. Numerous valuable analysis results have been developed for RFMs, we can only mention a few. In [27] it is shown that RFMs with a tube-like structure have a unique asymptotically stable equilibrium point within the invariant domain of their dynamics. RFMs with a ring topology are studied in [32] where it is shown that trajectories converge to equilibria within the compatibility classes of the state space defined by the initial conditions. The dynamics and stability of RFMs under periodic excitation is analyzed in [26]. We also mention that ODE models with essentially the same structure can be obtained by an appropriate finite volume discretization of

hyperbolic partial differential equations describing the flow of material or vehicles [25].

The structure of the paper is the following. In Sect. 2, we introduce the basic notions and known results for the ODE and Petri net representation of kinetic models. Section 3 describes the studied kinetic compartmental model class. Section 4 contains the persistence analysis results through the characterization of siphons, while the stability results are summarized in Sect. 5. Finally, the brief summary of the results is given in Sect. 6.

2 Background and notations

In this section, we introduce the class of kinetic systems and their representation in the form of Petri nets. Throughout the paper, we will use the following notations.

2.1 Kinetic systems

For the characterization of kinetic systems (also called chemical reaction networks or briefly, CRNs), we will use the notations used in [16], where more details can be found. A kinetic model contains M species denoted by $\mathcal{X} = \{X_1, \ldots, X_M\}$, and the corresponding species vector is given as $X = [X_1 \ \ldots \ X_M]^T$. Species are transformed into each other through *elementary reaction steps* of the form

$$C_j \to C_j', \quad j = 1, \ldots, R \tag{1}$$
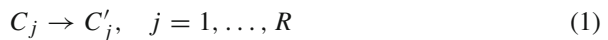
where $C_j = y_j^T X$ and $C_j' = y_j'^T X$ are the *complexes* with the *stoichiometric coefficient vectors* $y_j, y_j' \in \overline{\mathbb{Z}}_+^M$ for $j = 1, \ldots, R$. The transformation shown in Eq. (1) means that during an elementary reaction step between the *reactant complex* C_j and *product complex* C_j', $[y_j]_i$ molecules of species X_i are consumed, and $[y_j']_i$ molecules of X_i are produced for $i = 1, \ldots, M$. The reaction (1) is called an *input (output) reaction of species* X_i if $[y_j']_i > 0$ ($[y_j]_i > 0$).

The directed graph containing the complexes as vertices and reactions as directed edges is called the *reaction graph* of a CRN. A directed graph is *strongly connected* if there exists a directed path between any pair of its vertices in both directions. A *strong component* of a directed graph is a maximal strongly connected subgraph. A *weakly connected component* of a directed graph is a subgraph where all vertices are connected to each other by some (not necessarily directed) path. A reaction graph is called *weakly reversible* if each weakly connected component of it is a strong component. Weak reversibility is equivalent to the property that each directed edge (reaction) is a part of a directed cycle in the reaction graph.

Let $x(t) \in \overline{\mathbb{R}}_+^M$ denote the state vector corresponding to X for any $t \geq 0$ (in a chemical context, x is the vector of concentrations of the species in X). Then the ODEs describing the evolution of x in the kinetic system containing the reactions (1)

are given by

$$\dot{x} = \sum_{i=1}^{R} \mathcal{K}_i(x)[y_i' - y_i], \quad x(0) \in \overline{\mathbb{R}}_+^M \tag{2}$$

where $\mathcal{K}_i : \overline{\mathbb{R}}_+^M \longrightarrow \overline{\mathbb{R}}_+$ is the *rate function* corresponding to reaction step i, determining the velocity of the transformation [16]. For the rate functions, we assume the following for $i = 1, \ldots, R$:

(A1) $\mathcal{K}_i$ is differentiable,
(A2) $\dfrac{\partial \mathcal{K}_i(x)}{\partial x_j} \geq 0$ if $[y_i]_j > 0$, and $\dfrac{\partial \mathcal{K}_i(x)}{\partial x_j} = 0$ if $[y_i]_j = 0$,
(A3) $\mathcal{K}_i(x) = 0$ whenever $x_j = 0$ such that $j \in \mathrm{supp}(y_i)$.

The above properties guarantee the local existence and uniqueness of the solutions as well as the invariance of the nonnegative orthant for the dynamics in Eq. (2). From now on, a reaction from complex C_i to complex C_i' with rate function $\mathcal{K}_i$ will be denoted as

$$C_i \xrightarrow{\mathcal{K}_i} C_i' \tag{3}$$

The dynamics of a kinetic system (2) is called *persistent* if no trajectory that starts in the positive orthant has an omega-limit point on the boundary of $\mathbb{R}_+^M$.

A set of nonlinear ODEs given as $\dot{x} = f(x)$ is called *kinetic* if it can be written in the form (2) with appropriate rate functions $\mathcal{K}_i$. We remark that the representation (2) of a kinetic ODE is generally non-unique even if the rate functions are polynomial, and assumed to be fixed [1].

An important special case in the theory of CRNs is *mass action kinetics* when the rate function is given in the following monomial form
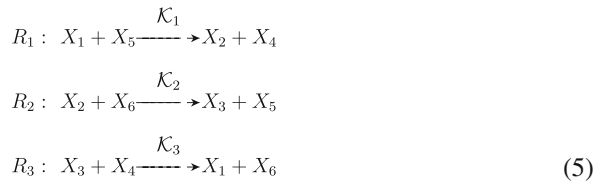
$$\mathcal{K}_i(x) = k_i \prod_{j=1}^{M} x_j^{[y_i]_j}, \quad i = 1, \ldots, R \tag{4}$$

where $k_i > 0$ for $i = 1, \ldots, R$ are the *reaction rate coefficients*.

A *positive linear conserved quantity* (or positive linear first integral) for a CRN is defined as $c^T x$ for which $c^T \dot{x}(t) = 0$ for $t \geq 0$, where $c \in \overline{\mathbb{R}}_+^M$ and $c \neq 0$. We say that a set of species $\{X_{i_1}, \ldots, X_{i_k}\} \subseteq \mathcal{X}$ *defines a positive linear conserved quantity* if there exists $c \in \overline{\mathbb{R}}_+^k$ for which $\sum_{j=1}^{k} c_j \dot{x}_{i_j}(t) = 0$ for $t \geq 0$.

Example 1

Consider the following CRN given by $\mathcal{X} = \{X_1, \ldots, X_6\}$ and the reactions

$$R_1: \quad X_1 + X_5 \xrightarrow{\mathcal{K}_1} X_2 + X_4$$

$$R_2: \quad X_2 + X_6 \xrightarrow{\mathcal{K}_2} X_3 + X_5$$

$$R_3: \quad X_3 + X_4 \xrightarrow{\mathcal{K}_3} X_1 + X_6 \tag{5}$$

Furthermore, assume that the reaction rates obey the mass action kinetics described in Eq. (4), i.e.

$$\mathcal{K}_1(x) = k_1 x_1 x_2, \quad \mathcal{K}_2(x) = k_2 x_2 x_6, \quad \mathcal{K}_3(x) = k_3 x_3 x_4 \tag{6}$$

Then according to Eq. (2), the ODEs of the system can be written as

$$\begin{aligned}
\dot{x}_1 &= -k_1 x_1 x_5 + k_3 x_3 x_4 \\
\dot{x}_2 &= k_1 x_1 x_5 - k_2 x_2 x_6 \\
\dot{x}_3 &= k_2 x_2 x_6 - k_3 x_3 x_4 \\
\dot{x}_4 &= k_1 x_1 x_5 - k_3 x_3 x_4 \\
\dot{x}_5 &= -k_1 x_1 x_5 + k_2 x_2 x_6 \\
\dot{x}_6 &= -k_2 x_2 x_6 + k_3 x_3 x_4
\end{aligned} \tag{7}$$

It can be checked from (7) that $\sum_{i=1}^{6} x_i$ is a positive linear conserved quantity for the kinetic system (5).

2.2 Petri net representation of CRNs and persistence conditions

If we consider each molecule and reaction individually, CRNs can be described in the framework of discrete event systems [30], and modeled e.g., by Petri nets [7]. Moreover, certain properties of the corresponding Petri net have fundamental consequences on the continuous dynamics of the studied CRN.

A *Petri net* is a directed bipartite graph $G = (V, E)$, where $V = \{v_1, v_2, \ldots, v_n\}$ is a set of vertices and $E = \{e_1, \ldots, e_k\}$ is a set of directed edges, i.e. $e_i = (v_j, v_k)$, where $v_j, v_k \in V$. The set of vertices can be partitioned into two disjoint sets, the set of places denoted by $P = \{p_1, \ldots, p_{n_p}\}$, and $T = \{t_1, \ldots, t_{n_t}\}$ which is the set of transitions, where $P \cup T = V$, and $P \cap T = \emptyset$. Moreover, for any $e_i = (v_j, v_k) \in E$, either $v_j \in P$ and $v_k \in T$ or vice versa.

The state of a Petri net is given by the number of *tokens* assigned to places. This can be characterized by a *marking* $\mu : P \longrightarrow \mathbb{N}_0$. Obviously, a marking can be given as an integer vector of size $|P|$. The places from which edges point to a transition are called the *input places of the transition*, while the places to which edges run from

Fig. 1 Petri net representation of the CRN described in Example 1. All edge weights are equal to 1
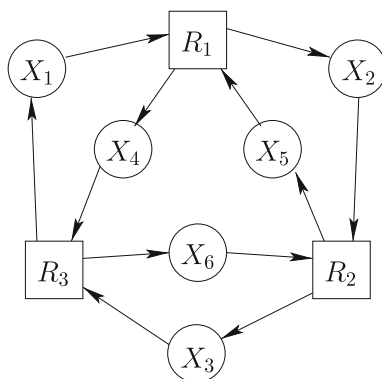
a transition are the *output places of the transition*. The input and output places of a transition t_i are denoted by $\text{In}(t_i)$ and $\text{Out}(t_i)$, respectively. Analogously, we can define the *input* and *output transitions* of a place p_i denoted by $\text{In}(p_i)$ and $\text{Out}(p_i)$, respectively. Positive integers are assigned to each directed edge through the *weighting* $W : (P \times T) \cup (T \times P) \longrightarrow \mathbb{N}$. A transition t_j is *enabled* if there are enough tokens in each of its input places: i.e., if $\forall \ p_i \in \text{In}(t_j)$: $\mu(p_i) \geq W(p_i, t_j)$. During the *firing* of an enabled transition t_j, $W(p_i, t_j)$ tokens are consumed from each $p_i \in \text{In}(t_j)$, and $W(t_j, p_k)$ tokens are added to each $p_k \in \text{Out}(t_j)$.

The dynamical behaviour of a Petri net is characterized by the sequence of transitions from an initial marking μ_0. Obviously, several transitions may be enabled at the same time which can fire in any order. Therefore, the execution (simulation) of Petri nets is generally nondeterministic.

It can be seen from the above, that Petri nets can be assigned in a straightforward way to kinetic systems. In such a modeling framework, places correspond to species, and transitions represent reactions. The input and output places of a transition correspond to the species of the reactant and the product complexes, respectively. The weights of the input and output edges of a transition are the stoichiometric coefficients of the species of the reactant and product complexes, respectively. For each place, the number of tokens show the actual number of molecules of the corresponding species. The Petri net representation of the CRN in Example 1 is shown in Fig. 1, where species (places) and reactions (transitions) are denoted by circles and rectangles, respectively.

A non-empty set of places $\sigma \subset P$ is called a *siphon* if each input transition associated to σ is also an output transition associated to σ. A siphon is *minimal* if it does not contain (strictly) any other siphons. Naturally, the union of siphons is a siphon, too. With some abuse of notation, a set of species Σ in a CRN will also be called a siphon if the places associated to the species of Σ form a siphon in the Petri net of the reaction network.

We will use a fundamental result from [7] which can be re-written as follows.

Theorem 2.1 *(Sufficient persistence conditions from [7]) The dynamics of a CRN of the form (2) is persistent if*

(1) *There exists a positive linear conserved quantity $c^T x$ for the dynamics, where $c \in \mathbb{R}_+^n$.*

(2) *Each siphon of the CRN contains a subset of species which define a positive linear conserved quantity for the dynamics.*

The practical difficulty in applying Theorem 1 is that the number of siphons generally grows exponentially with the network size [36], although there exist several computational approaches for the enumeration of all (minimal) siphons [19].

Let us revisit Example 1 to illustrate the conditions of Theorem 2.1. It is easy to see that condition (1) is fulfilled, since $I_0 = \sum_{i=1}^{6} x_i$ is a positive conserved quantity for the system as it was written in Sect. 2.1. For condition (2), it can be checked from Fig. 1 that five minimal siphons exist in the Petri net, namely $\Sigma_1 = \{X_1, X_2, X_3\}$, $\Sigma_2 = \{X_4, X_5, X_6\}$, $\Sigma_3 = \{X_1, X_4\}$, $\Sigma_4 = \{X_2, X_5\}$, and $\Sigma_5 = \{X_3, X_6\}$. Since $I_1 = x_1 + x_2 + x_3$, $I_2 = x_4 + x_5 + x_6$, $I_3 = x_1 + x_4$, $I_4 = x_2 + x_5$, and $I_5 = x_3 + x_6$ are also positive linear first integrals containing the state variables of Σ_1, Σ_2, Σ_3, Σ_4, and Σ_5, respectively, condition (2) is fulfilled, too.

3 The studied compartmental model class

In this paper, we consider a subclass of flow models equipped with a network structure. In such models, we have interconnected compartments and items (e.g., molecules, particles, or vehicles) moving between them. The compartments have finite capacities, i.e. we assume that there are well-defined upper limits for the number of items placed in the compartments at any time instant. The transition rate of items between two compartments depend on the number of particles in the source compartment and on the amount of available space in the target compartment.

3.1 Directed graph of the compartmental structure

The structure of a compartmental model showing the possible directions of flows between the compartments can be described by a directed graph as follows.

Definition 3.1 The directed graph $D = (Q, A)$ called **compartmental graph** describes the structure of the compartmental model, where the set $Q = \{q_1, \ldots, q_m\}$ of vertices correspond to the compartments, and the possible transitions are represented by directed edges of the set $A \subseteq Q \times Q$. The directed edge $a_{ij} := (q_i, q_j) \in A$ represents the transition from the compartment q_i into q_j.

Naturally, loop edges are not allowed in the compartmental graph, since the immediate transition from a compartment into itself does not induce any change. Furthermore, multiple identically directed edges are also not allowed between two compartments. If there exists a directed edge (q_i, q_j) in the compartmental graph, then q_i is called the *donor* of compartment q_j, and q_j is the *recipient* of compartment q_i.

The directed graph of a strongly connected triangular compartmental model with $Q = \{q_1, q_2, q_3\}$, and $A = \{(q_1, q_2), (q_2, q_3), (q_3, q_1)\} = \{a_{12}, a_{23}, a_{31}\}$ is shown in Fig. 2.

Fig. 2 Directed graph of a simple triangular compartmental model
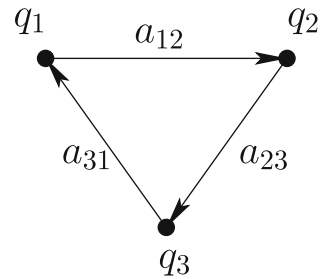
3.2 Kinetic representation of compartmental models

We assign a CRN to a compartmental model $D = (Q, A)$ containing m compartments as follows. The set of species is $\Sigma = \mathcal{N} \cup \mathcal{S}$, where $\mathcal{N} = \{N_1, \ldots N_m\}$ and $\mathcal{S} = \{S_1, \ldots S_m\}$, where N_i and S_i represent the amount of particles and the available space in compartment q_i, respectively. To each transition (directed edge) a_{ij} in D we assign the following reaction

$$N_i + S_j \xrightarrow{\mathcal{K}_{ij}} N_j + S_i \tag{8}$$

Eq. (8) shows that during an elementary step of the transition from compartment q_i to q_j, the amount of content (e.g., particles, material) in compartment q_i is decreased by one unit, and the number of particles in compartment q_j is increased by one unit. Parallelly, the amount of free space is increased in compartment q_i and decreased in compartment q_j. It is also visible that a necessary condition for any transition is that there is at least one particle in compartment q_i and at least one available space in q_j. The rate (velocity) of the transition is determined by the rate function $\mathcal{K}_{ij}$. Let us denote the continuous amount (or concentration) of particles and free space in compartment q_i by n_i and s_i, respectively. Moreover, let $\mathcal{D}_i$ and $\mathcal{R}_i$ denote the index sets of the donor and recipient compartments of q_i, respectively. Then, using Eq. (2), the dynamics of n_i and s_i can be written as

$$\begin{aligned} \dot{n}_i &= \sum_{j \in \mathcal{D}_i} \mathcal{K}_{ji}(n_j, s_i) - \sum_{j \in \mathcal{R}_i} \mathcal{K}_{ji}(n_i, s_j) \\ \dot{s}_i &= - \sum_{j \in \mathcal{D}_i} \mathcal{K}_{ji}(n_j, s_i) + \sum_{j \in \mathcal{R}_i} \mathcal{K}_{ji}(n_i, s_j) \end{aligned} \tag{9}$$

It is visible from Eq. (9) that $c_i := n_i + s_i$ is constant for any compartment q_i, therefore, c_i will be called the *capacity* of q_i for $i = 1, \ldots, m$. From the CRN defined by the species and reactions in Eq. (8), we can give the Petri net representation of a compartmental model as it is described in Sect. 2.2.

Consider the compartmental model shown in Fig. 2. The associated CRN model is the following

$$N_1 + S_2 \xrightarrow{\bar{\mathcal{K}}_{12}} N_2 + S_1$$

$$N_2 + S_3 \xrightarrow{\bar{\mathcal{K}}_{23}} N_3 + S_2$$

$$N_3 + S_1 \xrightarrow{\bar{\mathcal{K}}_{31}} N_1 + S_3 \tag{10}$$

It is easy to see that the CRN (10) is identical to (5) in Example 1 with $N_i = X_i$ and $S_i = X_{i+3}$ for $i = 1, 2, 3$, and $\bar{\mathcal{K}}_{12} = \mathcal{K}_1$, $\bar{\mathcal{K}}_{23} = \mathcal{K}_2$, $\bar{\mathcal{K}}_{31} = \mathcal{K}_3$. Therefore, the corresponding Petri net is the same as the one shown in Fig. 1. Note that both the compartmental graph and the Petri net of the model are strongly connected. However, as it is visible from the disjoint complexes of the reactions listed in Eq. (10), the corresponding reaction graph is not weakly reversible and therefore not strongly connected.

It is important to remark that the model class introduced in this section includes as special cases certain finite volume discretizations of hyperbolic partial differential equations applied e.g., in (traffic) flow modeling [23, 25], and popular kinetic models for the description of simultaneous mRNA translation and competition for ribosomes [31, 32].

4 Persistence analysis

Proposition 4.1 *If the compartmental graph D of a model is strongly connected then the corresponding Petri net $\mathcal{P}(D)$ as a directed graph is also strongly connected.*

A directed graph is strongly connected if for any two vertices v and w there is a directed path from v to w. The idea of the proof is that if there are two paths $P_1 = v_1 v_2 \ldots v_k$ and $P_2 = v_k v_{k+1} \ldots v_n$, then by concatenating P_2 after P_1 we get a walk from v_1 to v_n. It will be a walk, and not necessarily a path, since there might be identical vertices in the two concatenated paths. In this case by omitting the loop created between the first and last occurrences of a vertex we can get a shorter walk between the same endpoints. By a series of such steps the repeated occurences can be eliminated, and we get a path. Consequently, if there exists a walk from v to w, then there exists a path from v to w as well, and for a graph to be strongly connected, it is enough to show the existence of walks instead of paths between any two vertices.

Proof Recall from Sect. 3 that the Petri net corresponding to the CRN representation of a compartmental model contains two types of vertices, N_i and S_j representing the number of molecules in compartment q_i and the empty spaces in compartment q_j, respectively. The transition between the compartments q_i and q_j also corresponds to a vertex, but this is a different, reaction type vertex called R_{ij}.

If in the compartmental graph D there is a transition from compartment q_i to q_j, then in the Petri net $\mathcal{P}(D)$ there exists a corresponding reaction vertex R_{ij}, which is
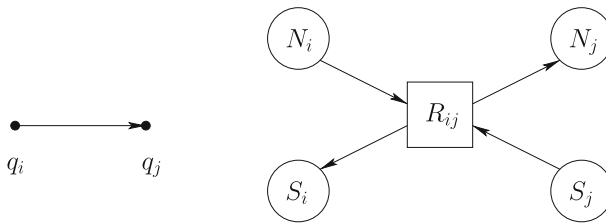
Fig. 3 Representations of the same transition in the compartmental graph D and in the Petri net $\mathcal{P}(D)$

incident to directed edges from vertices N_i and S_j, and to vertices N_j and S_i. It can be seen in Fig. 3 that a transition in the compartmental graph induces a path in the set $\mathcal{N}$ of the Petri net in the same direction, and in the set $\mathcal{S}$ in the opposite direction.

We will examine the existence of paths between different types of vertices of the Petri net separately, i.e. there will be four cases:

1. Path from N_i to N_j
   Since the compartmental graph D is strongly connected, it contains a directed path from q_i to q_j:

$$q_i \rightarrow q_{k_1} \rightarrow q_{k_2} \rightarrow \ldots \rightarrow q_{k_l} \rightarrow q_j$$

   This implies a directed path in the Petri net $\mathcal{P}(D)$ from N_i to N_j with twice the length of the path in the compartmental graph:

$$N_i \rightarrow R_{ik_1} \rightarrow N_{k_1} \rightarrow R_{k_1k_2} \rightarrow N_{k_2} \rightarrow \ldots \rightarrow N_{k_l} \rightarrow R_{k_lj} \rightarrow N_j$$

2. Path from S_i to S_j
   The existence of such a path can be proven similarly as in the previous case. Because of the strong connectivity of the compartmental graph D there exists a directed path from q_j to q_i:

$$q_j \rightarrow q_{m_1} \rightarrow q_{m_2} \rightarrow \ldots \rightarrow q_{m_p} \rightarrow q_i$$

   This implies a directed walk in the Petri net $\mathcal{P}(D)$ from S_i to S_j:

$$S_j \leftarrow R_{jm_1} \leftarrow S_{m_1} \leftarrow R_{m_1m_2} \leftarrow S_{m_2} \leftarrow \ldots \leftarrow S_{m_p} \leftarrow R_{m_pi} \leftarrow S_i$$

3. Path from N_i to S_j
   The strong connectivity of the compartmental graph D implies that for every vertex there is at least one directed edge starting there, consequently there must be a vertex q_k to which there is a transition from q_i. Through the reaction vertex in the Petri net there is a path $N_i \rightarrow R_{ik} \rightarrow S_i$. By concatenating this path with the existing path from S_i to S_j we get a walk from N_i to S_j.
4. Path from S_i to N_j
   The strong connectivity of the compartmental graph D also implies that for every vertex there is at least one directed edge going there, consequently there must be a
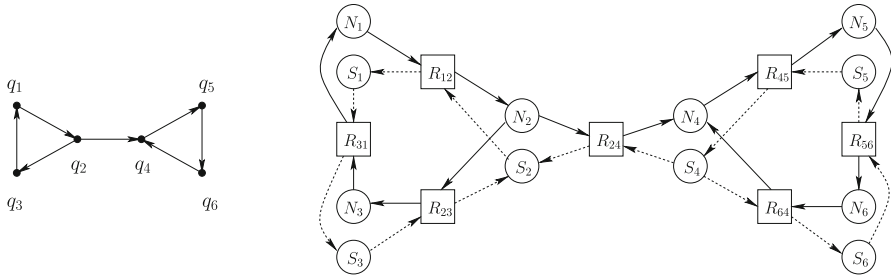
Fig. 4 Strongly connected Petri net $\mathcal{P}(D)$ corresponding to a not strongly connected compartmental graph D

vertex q_l from which there is a transition to q_i. It follows that in the Petri net there is a path $S_i \rightarrow R_{li} \rightarrow N_i$. By concatenating this path with the existing path from N_i to N_j we get a walk from S_i to N_j.

It can be seen that since there exist paths connecting the type $\mathcal{N}$ and type $\mathcal{S}$ vertices, then there are paths between the reaction type vertices as well. $\qquad \square$

Remark 4.2 The implication in the other direction is not true. It is possible that the Petri net is strongly connected but the corresponding compartmental graph is not. Such an example can be seen in Fig. 4.

In the following part of this section we will examine the structure of siphons in the Petri net. According to the definition it is a subset of $\Sigma = \mathcal{N} \cup \mathcal{S}$ for which every input reaction is also an output reaction.

Proposition 4.3 *If the compartmental graph D is strongly connected and $\mathcal{N}' \subseteq \mathcal{N}$ is a siphon in the Petri net $\mathcal{P}(D)$, then $\mathcal{N}' = \mathcal{N}$ must hold.*

In other words, if a siphon of the Petri net contains vertices only from the set $\mathcal{N}$, then it contains all of them.

Proof Let us assume by contradiction that $\mathcal{N}'$ is a siphon in the Petri net, which is a real subset of the vertex set $\mathcal{N}$.

This set corresponds to the set $Q' = \{q_i \mid N_i \in \mathcal{N}'\}$, which is a real subset of the vertex set Q of the compartmental graph D, i.e. $Q' \neq \emptyset$ and $Q \setminus Q' \neq \emptyset$ hold. Since the compartmental graph D is strongly connected, there must be vertices $q_i \in Q'$ and $q_j \in Q \setminus Q'$ so that there is a directed edge in D from q_j to q_i.

This edge represents the transition corresponding to the reaction R_{ji} in the Petri net, which connects the vertices N_j, S_j, N_i and S_i. Since the vertex N_i is in the siphon $\mathcal{N}'$, the reaction R_{ji} is an input reaction of the siphon $\mathcal{N}'$, so by the definition of siphons, R_{ji} must be an output reaction to $\mathcal{N}'$ as well. For this to hold S_i or N_j should be in the set $\mathcal{N}'$. However, $S_i \notin \mathcal{N}'$ since by its definition $\mathcal{N}'$ contains only type $\mathcal{N}$ vertices, and $N_j \notin \mathcal{N}'$ since $q_j \notin Q'$. The original assumption leads to contradiction, meaning that the siphon $\mathcal{N}'$ cannot be a real subset of $\mathcal{N}$.

However, $\mathcal{N}$ itself is a siphon, since every reaction R_{kl} is an input reaction to the vertex N_l and an output reaction from the vertex N_k. $\qquad \square$

A similar property is true for the species set $\mathcal{S}$, and the proof is based on the same idea.

Proposition 4.4 *If the compartmental graph D is strongly connected and $\mathcal{S}' \subseteq \mathcal{S}$ is a siphon in the Petri net $\mathcal{P}(D)$, then $\mathcal{S}' = \mathcal{S}$ must hold.*

If a siphon in the Petri net contains both types of vertices, a different type of structural property can be formulated.

Proposition 4.5 *If the compartmental graph D is strongly connected and $\mathcal{T} \subseteq \mathcal{N} \cup \mathcal{S}$ is a siphon in the corresponding Petri net $\mathcal{P}(D)$ for which $\mathcal{T} \cap \mathcal{N} \neq \emptyset$ and $\mathcal{T} \cap \mathcal{S} \neq \emptyset$, then there is an index $i \in \{1, 2, \ldots, m\}$ for which $N_i \in \mathcal{T}$ and $S_i \in \mathcal{T}$.*

In other words, if a siphon in the Petri net $\mathcal{P}(D)$ of a strongly connected compartmental graph D contains type $\mathcal{N}$ and type $\mathcal{S}$ vertices as well, then there is a vertex q_i in the compartmental graph D for which the corresponding vertices N_i and S_i are both contained by the siphon.

Proof Let us assume by contradiction that there is no such index. In this case the siphon $\mathcal{T}$ is of the form $\mathcal{N}' \cup \mathcal{S}'$, where the corresponding vertex subsets $Q_1 = \{q_i \mid N_i \in \mathcal{N}'\}$ and $Q_2 = \{q_j \mid S_j \in \mathcal{S}'\}$ of the compartmental graph D are disjoint. Consequently, the sets Q_1, Q_2 and $Q \setminus (Q_1 \cup Q_2)$ form a partition of the vertices of D. (The set $Q \setminus (Q_1 \cup Q_2)$ might be empty, in this case we have a partition with two class instead of three.)

Since the compartmental graph D is strongly connected, there must be an edge going into the set Q_1 from either or both of the other sets, i.e. there exist vertices $q_j \in Q_1$ and $q_i \in Q \setminus Q_1$ so that $q_i q_j$ is a directed edge in D. The corresponding reaction R_{ij} in the Petri net $\mathcal{P}(D)$ is an input reaction to N_j in the siphon $\mathcal{T}$, therefore it must be an output reaction as well. For this to be fulfilled, S_j or N_i should be in the siphon $\mathcal{T}$. By the assumption $\mathcal{T}$ cannot contain S_j since $N_j \in \mathcal{T}$, and $\mathcal{T}$ cannot contain N_i since the corresponding vertex q_i is not in the set Q_1. This is a contradiction, consequently there must be an index i for which $N_i \in \mathcal{T}$ and $S_i \in \mathcal{T}$ hold. □

Corollary 4.6 *A siphon in the Petri net of a strongly connected compartmental graph either contains the N_i and S_i vertices corresponding to the same compartment q_i, or it contains all the vertices of the same type $\mathcal{N}$ or $\mathcal{S}$.*

Proposition 4.7 *In the Petri net $\mathcal{P}(D')$ of any compartmental graph D' the set $\{N_i, S_i\}$ is a siphon for all $i \in \{1, 2, \ldots, m\}$.*

Proof If R is an input reaction to the set $\{N_i, S_i\}$, then it corresponds to a transition to or from the compartment q_i.

If R is an input reaction to the vertex N_i in the Petri net $\mathcal{P}(D')$, then it represents a transition a_{ji} from some compartment q_j to the compartment q_i in the compartmental graph D'. In this case $R = R_{ji}$, and in the Petri net $\mathcal{P}(D')$ this reaction is also an output reaction from the vertex S_i.

If R is an input reaction to the vertex S_i in the Petri net $\mathcal{P}(D')$, then it represents a transition a_{ik} from the compartment q_i to some compartment q_k in the compartmental graph D'. In this case $R = R_{ik}$, and in the Petri net $\mathcal{P}(D')$ this reaction is also an output reaction from the vertex N_i (Fig. 5). □
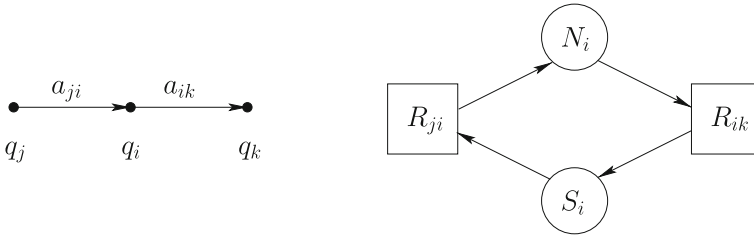
Fig. 5 The vertex set $\{N_i, S_i\}$ is a siphon in the Petri net $\mathcal{P}(D')$

Corollary 4.8 *Among the siphons contained by the Petri net of a strongly connected compartmental graph the ones that are minimal with respect to containment are only the sets $\mathcal{N}$, $\mathcal{S}$ and the sets $\{N_i, S_i\}$ for all indices $i \in \{1, 2, \ldots, m\}$.*

Proof Let $\mathcal{T}$ refer to a siphon in the Petri net of a strongly connected compartmental graph.

If the siphon $\mathcal{T}$ contains just type $\mathcal{N}$ or just type $\mathcal{S}$ vertices, then by Propositions 4.3 and 4.4 it must contain all the vertices of that type. Consequently, the sets $\mathcal{N}$ and $\mathcal{S}$ are siphons that are minimal with respect to containment.

If the siphon $\mathcal{T}$ contains both type $\mathcal{N}$ and type $\mathcal{S}$ vertices, then by Proposition 4.5 there must be an index $i \in \{1, 2, \ldots, m\}$ for which $\{N_i, S_i\} \subseteq \mathcal{T}$ holds. In Proposition 4.7 it was proven that the set $\{N_i, S_i\}$ for all indices is a siphon in the Petri net of a strongly connected compartmental graph. Consequently, if $\mathcal{T}$ is not equal to $\{N_i, S_i\}$ for some index i, then it is not minimal with respect to containment.

A vertex N_i in itself cannot be a siphon in the Petri net of a strongly connected compartmental graph, it follows from Proposition 4.3. Similarly, S_i in itself cannot be a siphon there, based on Proposition 4.4. Consequently, for every index $i \in \{1, 2, \ldots, m\}$ the set $\{N_i, S_i\}$ is a minimal siphon with respect to containment.

$\square$

Corollary 4.9 *The dynamics given in Eq. (9) is persistent if the corresponding compartmental graph $D = (Q, A)$ is strongly connected.*

Proof The statement is an immediate consequence of Theorem 2.1 and Corollary 4.6.

$\square$

5 Stability results

Consider a system given by (9) with strongly connected compartmental structure. Using the fact that for each q_i compartment $n_i = c_i - s_i$ we can rewrite the system in the reduced state-space as

$$\dot{n}_i = \sum_{j \in \mathcal{D}_i} \mathcal{K}_{ji}(n_j, c_i - n_i) - \sum_{j \in \mathcal{R}_i} \mathcal{K}_{ij}(n_i, c_j - n_j). \tag{11}$$

The state-space of the above system is $C := [0, c_1] \times [0, c_2] \times \cdots \times [0, c_m]$ and let ∂C denote the boundary of C; that is, $\partial C = C \backslash \text{int}(C)$.

The Jacobian of (11) is given by

$$
[J(n)]_{ik} = \begin{cases}
-\sum_{j \in \mathcal{D}_i} \frac{\partial \mathcal{K}_{ji}(n_j, c_i - n_i)}{\partial n_i} - \sum_{j \in \mathcal{R}_i} \frac{\partial \mathcal{K}_{ij}(n_i, c_j - n_j)}{\partial n_i} & \text{if } i = k, \\
\frac{\partial \mathcal{K}_{ki}(n_k, c_i - n_i)}{\partial n_k} & \text{if } k \in \mathcal{D}_i \text{ and } k \notin \mathcal{R}_i, \\
\frac{\partial \mathcal{K}_{ik}(n_i, c_k - n_k)}{\partial n_k} & \text{if } k \notin \mathcal{D}_i \text{ and } k \in \mathcal{R}_i, \\
\frac{\partial \mathcal{K}_{ki}(n_k, c_i - n_i)}{\partial n_k} + \frac{\partial \mathcal{K}_{ik}(n_i, c_k - n_k)}{\partial n_k} & \text{if } k \in \mathcal{D}_i \text{ and } k \in \mathcal{R}_i, \\
0 & \text{otherwise.}
\end{cases} \tag{12}
$$

The (A2) property of the rate functions imply that each diagonal entry is nonpositive and each off-diagonal entry is nonnegative. Since the sum of each column is zero, we conclude that the system is compartmental in the sense of [22]. Systems satisfying the latter property are also called cooperative.

The following lemmata and proofs will adapt the ideas of [27] and [32] for the studied more general system class. Moreover, we will also use the persistence result of Corollary 4.9.

Lemma 5.1 *Consider a compartmental system of the form (11) with a strongly connected compartmental structure. Then, for any $n(0) \in \text{int}(C)$ the solution satisfies $n(t) \in \text{int}(C)$ for any $t \geq 0$.*

In other words, $\text{int}(C)$ is an invariant set of such a system.

Proof To obtain a contradiction, suppose that there exists a (minimal) time $\tau > 0$ such that $n(\tau) \in \partial C$. We need to consider the following two cases.

1. There exists an empty compartment. In this case, due to the strongly connected structure, there must exist an empty compartment with at least one non-empty donor compartment as well. To see this, consider a directed path from any non-empty compartment to any empty compartment. Stepping backwards from the empty compartment along this path until we reach a non-empty compartment establishes our assertion.

   Let i be an index such that $n_i(\tau) = 0$ and $n_k(\tau) > 0$ holds for some $k \in \mathcal{D}_i$. Then (11) takes the form

$$
\dot{n}_i(\tau) = \sum_{j \in \mathcal{D}_i} \mathcal{K}_{ji}(n_j, c_i) \geq \mathcal{K}_{ki}(n_k, c_i) > 0 \tag{13}
$$

   which means that $\dot{n}_i(t) > 0$ on the interval $[\tau - \sigma, \tau]$ for some $\sigma > 0$. This leads to a contradiction with $n_i(\tau) = 0$, further implying that there are no empty compartments altogether.

2. There exists a full compartment. In this case, by a similar argument, there must exist a full compartment with at least one non-full recipient compartment as well; that is, there exists an index i such that $n_i(\tau) = c_i$ and $n_k(\tau) < c_k$ holds for some

$k \in \mathcal{R}_i$. Then (11) takes the form

$$\dot{n}_i(\tau) = -\sum_{j \in \mathcal{R}_i} \mathcal{K}_{ij}(c_i, c_j - n_j) \leq -\mathcal{K}_{ik}(c_i, c_k - n_k) < 0 \tag{14}$$

which means that $\dot{n}_i(t) < 0$ on the interval $[\tau - \sigma, \tau]$ for some $\sigma > 0$. This leads to a contradiction with $n_i(\tau) = c_i$, further implying that there are no full compartments altogether.

□

Let $0^{(m)}, c^{(m)} \in \mathbb{R}^m$ be defined by

$$0^{(m)} = \begin{bmatrix} 0 \\ 0 \\ \vdots \\ 0 \end{bmatrix} \qquad c^{(m)} = \begin{bmatrix} c_1 \\ c_2 \\ \vdots \\ c_m \end{bmatrix}.$$

Lemma 5.2 *Consider a compartmental system of the form (11) with a strongly connected compartmental structure. Then, for any $n(0) \in \partial C$, $n(0) \neq 0^{(m)}$, $n(0) \neq c^{(m)}$ the solution satisfies $n(\tau) \in \mathrm{int}(C)$ for some $\tau > 0$.*

Proof First we define the following boundary-repelling property.

(**BR**) For each $\delta > 0$ and sufficiently small $\Delta > 0$, there exists $K = K(\delta, \Delta) > 0$ such that for each $t \geq 0$

1. the conditions

   (a) $n_i(t) \leq \Delta$,
   (b) there exists $k \in \mathcal{D}_i$ such that $n_k(t) \geq \delta$

   imply $\dot{n}_i(t) \geq K$, and
2. the conditions

   (a) $n_i(t) \geq c_i - \Delta$
   (b) there exists $k \in \mathcal{R}_i$ such that $n_k(t) \leq c_k - \delta$

   imply $\dot{n}_i(t) \leq -K$.

(11) satisfies the above property. To see this, consider any compartment q_i. Without the loss of generality we can assume that $\mathcal{D}_i$ contains at least one index, let this be k. In this case

$$\dot{n}_i(t) \geq \mathcal{K}_{ki}(\delta, c_i - \Delta) - \sum_{j \in \mathcal{R}_i} \mathcal{K}_{ij}(\Delta, c_j) := K_1. \tag{15}$$

Similarly, we can assume that $\mathcal{R}_i$ contains at least one index, let this be l. In this case

$$\dot{n}_i(t) \leq \sum_{j \in \mathcal{D}_i} \mathcal{K}_{ji}(c_j, \Delta) - \mathcal{K}_{il}(c_i - \Delta, c_l - \delta) := -K_2. \tag{16}$$

The properties of the rate functions imply that for a sufficiently small Δ we have $K_1 > 0$ and $-K_2 < 0$, thus taking $K = \min\{K_1, K_2\}$ concludes our assertion.

Next, we will show that for each compartment $n_i(\tau) > 0$ holds for some $\tau > 0$.

Without the loss of generality we can assume that there exists an index i such that $n_i(t) \geq \epsilon_0$ on the interval $[0, \tau]$ for some $\epsilon_0 > 0$ and $\tau > 0$. Define $\tau_m = \frac{\tau}{m}$ and proceed by induction. For $k = 1, 2, \ldots, m$ we will define an appropriate $\epsilon_k > 0$ and show that the kth generation recipients of the compartment q_i have particle concentration of at least ϵ_k on the interval $[k\tau_m, \tau]$.

Pick any $j \in \mathcal{R}_i$ (first generation recipient) and sufficiently small $\Delta > 0$, define $K = K(\epsilon_0, \Delta)$ and $\epsilon_1 = \min\{\Delta, K\tau_m\}$ and let $t_0 \in [0, \tau_m]$ such that $n_j(t_0) \geq \epsilon_1$. Such a t_0 must exist, since assuming $n_j(t) < \epsilon_1 \leq \Delta$ for each $t \in [0, \tau_m]$ would imply via (**BR**) that $\dot{n}_j(t) \geq K$ for each $t \in [0, \tau_m]$. This further implies that $n_j(\tau_m) \geq n_j(0) + K\tau_m \geq \epsilon_1$. This leads to a contradiction with $n_j(\tau_m) < \epsilon_1$.

Our next claim is that $n_j(t) \geq \epsilon_1$ for each $t \in [t_0, \tau]$ and in particular $[\tau_m, \tau]$. Conversely, suppose that there exists some $t_1 \in (t_0, \tau]$ such that $\xi := n_j(t_1) < \epsilon_1$ and define $\sigma = \min\{t \in (t_0, \tau) : n_j(t) \leq \xi\}$. Since $n_j(\sigma) \leq \xi < \epsilon_1 \leq \Delta$, (**BR**) shows that $\dot{n}_j(\sigma) \geq K$; that is, $\dot{n}_j(t) > 0$ on the interval $[\sigma - \nu, \sigma]$ for some $\nu > 0$. But this would imply that $n_j(\sigma - \nu) < n_j(\sigma)$, contradicting the minimality of σ.

Define $K = K(\epsilon_1, \Delta)$ and $\epsilon_2 = \min\{\Delta, K\tau_m\}$ and repeat the above steps for the set $\mathcal{R}_j$ for $j \in \mathcal{R}_i$ (second generation recipients). In subsequent induction steps define $K = K(\epsilon_k, \Delta)$ and $\epsilon_{k+1} = \min\{\Delta, K\tau_m\}$ and repeat the above for the kth generation recipients of the compartment q_i. Since the compartments are strongly connected after at most m induction steps we conclude that $n_i(\tau) > 0$ for each $i = 1, 2, \ldots, m$.

To show that $n_i(\tau) < c_i$ holds as well, consider the complementary system obtained by rewriting (9) using $s_i = c_i - n_i$ as

$$\dot{s}_i = -\sum_{j \in \mathcal{D}_i} \mathcal{K}_{ji}(c_j - s_j, s_i) + \sum_{j \in \mathcal{R}_i} \mathcal{K}_{ij}(c_i - s_i, s_j). \tag{17}$$

Repeating the above steps for (17) shows that $s_i(\tau) > 0$, further implying that $n_i(\tau) < c_i$; that is, indeed $n(\tau) \in \mathrm{int}(C)$. □

Remark 5.3 The proof also shows that for each $\tau > 0$ there exists $\epsilon(\tau) > 0$ with $\epsilon(\tau) \to 0$ as $\tau \to 0$, such that $n(\tau) \in [\epsilon, c_1 - \epsilon] \times [\epsilon, c_2 - \epsilon] \times \cdots \times [\epsilon, c_m - \epsilon]$; that is, even if the initial value is on ∂C the orbit enters $\mathrm{int}(C)$ after an arbitrarily short time.

Remark 5.4 A similar argument shows that ∂C only contains the two trivial equilibria corresponding to an empty and a full network.

To see this, let us first assume that n^* is an equilibrium and for a compartment q_i we have $n_i^* = 0$. Then, by (11)

$$\dot{n}_i^* = \sum_{j \in \mathcal{D}_i} \mathcal{K}_{ji}(n_j^*, c_i) = 0$$

which is only possible if $n_j^* = 0$ for each $j \in \mathcal{D}_i$. Induction shows that $n^* = 0^{(m)}$.

Next, let us assume that for a compartment q_i we have $n_i^* = c_i$. Then, by (11)

$$\dot{n}_i^* = -\sum_{j \in \mathcal{R}_i} \mathcal{K}_{ij}(c_i, c_j - n_j^*) = 0$$

which is only possible if $n_j^* = c_j$ for each $j \in \mathcal{R}_i$. Induction shows that $n^* = c^{(m)}$.

For a given initial condition $a \in C$, let $\varrho(t, a)$ denote the solution at time t with $\varrho(0, a) = a$; that is $\varrho(t, a) = n(t)$ with $n(0) = a$. Since the total number of particles is conserved, the function $I : \mathbb{R}^m \mapsto \mathbb{R}$ defined by $I(y) = \sum_{i=1}^m y_i$ is a first integral. For $s \in \left[0, I(c^{(m)})\right]$ let $L_s \subset C$ be the level set of I; that is,

$$L_s = \{a \in C : I(a) = s\}. \tag{18}$$

Using the terminology of CRN theory [16], the level sets defined in Eq. (18) are also called *stoichiometric compatibility classes*.

Proposition 5.5 *Consider a compartmental system of the form* (11) *with a strongly connected compartmental structure. Then, for any* $s \in \left[0, I(c^{(m)})\right]$ *the set* L_s *contains a unique steady state* e_s *satisfying* $\lim_{t \to \infty} \varrho(t, a) = e_s$ *for any* $a \in L_s$.

Proof Since $L_0 = \{0^{(m)}\}$ and $\varrho(t, 0^{(m)}) = 0^{(m)}$, the statement holds for an empty network with $e_0 = 0^{(m)}$. Similarly, since $L_{I(c^{(m)})} = \{c^{(m)}\}$ and $\varrho(t, c^{(m)}) = c^{(m)}$, the statement holds for a full network with $e_{I(c^{(m)})} = c^{(m)}$.

Choose $s \in \left(0, I(c^{(m)})\right)$ and $a \in L_s$. By the strongly connected compartmental structure the Jacobian $J(n)$ is irreducible on int(C) but may become reducible on ∂C. However, Lemmas 5.1 and 5.2 along with Remark 5.3 show that (11) has repelling boundary; that is, $\varrho(t, a) \in$ int(C) after an arbitrarily short time even if $a \in L_s \cap \partial C$. As a consequence, (11) is a cooperative irreducible system evolving in int(C) admitting a first integral with positive gradient. The result [28,Theorem 10.] shows that L_s either has precisely one equilibrium that attracts the whole level set or has zero equilibria and each ω-limit set of the level set is empty. However, by the boundedness of the sequence $\{\varrho(k, a) : k = 1, 2, \ldots\} \subset$ int(C) the Bolzano-Weierstrass theorem implies that there is a convergent subsequence; that is, the ω-limit set of a cannot be empty. Furthermore, Corollary 4.9 implies that $\omega(a) \cap \partial C = \emptyset$ and the proof is complete. □

In the proofs above we used the notion of cooperative systems directly, however, the underlying theory involves so-called (strongly) monotone systems, which in our case, is a direct consequence of cooperativity, as shown by our next result.

For two points $x, y \in \mathbb{R}^m$, let

$$x \le y \quad \text{if } x_i \le y_i \text{ for } i = 1, 2, \ldots, m, \tag{19}$$

$$x < y \quad \text{if } x \le y \text{ and } x \ne y, \tag{20}$$

$$x \ll y \quad \text{if } x_i < y_i \text{ for } i = 1, 2, \ldots, m. \tag{21}$$

Proposition 5.6 *Consider a compartmental system of the form* (11) *with a strongly connected compartmental structure. Then, for any* $s \in \left[0, I(c^{(m)})\right]$ *and* $a, b \in L_s$, *the relation* $a \leq b$ *implies* $\varrho(t, a) \leq \varrho(t, b)$ *and* $a < b$ *implies* $\varrho(t, a) \ll \varrho(t, b)$ *for any* $t > 0$.

Proof If x or y is equal to $0^{(m)}$ or $c^{(m)}$, then the statement trivially holds. In any other case, use the proof of Proposition 5.5 to conclude that (11) is a cooperative irreducible system evolving in a convex and open set, namely, $\text{int}(C)$. The statement is a direct consequence of [28,Theorem 1., Theorem 3.]. □

Our final result in this topic gives further insight into the qualitative behaviour of (11).

Proposition 5.7 *Consider a compartmental system of the form* (11) *with a strongly connected compartmental structure. Then, for any* $a, b \in C$ *initial values and* $t \geq 0$

$$\|\varrho(t, a) - \varrho(t, b)\|_{\ell^1(\mathbb{R}^m)} \leq \|a - b\|_{\ell^1(\mathbb{R}^m)}.$$

In other words, using the usual $\ell^1(\mathbb{R}^m)$ norm, the distance of two trajectories at any given time cannot be larger than the distance of the initial values. In particular, if $b = e_{I(a)}$, then we find that the convergence to $e_{I(a)}$ is monotone.

Proof By [35,Chapter 2.2] the induced matrix measure by the ℓ^1 vector norm is

$$\mu(A) = \max_i \left\{ [A]_{ii} + \sum_{j \neq i} \left| [A]_{ji} \right| \right\}. \tag{22}$$

Since $\mu(J(n)) = 0$, the result [34,Theorem 1.] implies the assertion of the proposition. □

Remark 5.8 It is straightforward to extend our persistence and stability results to systems with a weakly reversible compartmental graph, when the dynamics unfold into isolated subsystems having strongly connected compartmental graphs. Furthermore, some of the above results on the qualitative behaviour, for example the monotonicity in Propositions 5.6 and 5.7 can be extended to systems with arbitrary compartmental topology.

6 Conclusions

The fundamental dynamical properties of a class of kinetic compartmental systems having finite capacities were studied in this paper. The transition of particles between compartments was described by a kinetic model. For persistence analysis, the Petri net representation of the CRN model was used. We showed that the Petri net of models having a strongly connected compartmental graph is also strongly connected. It was also shown that the dynamics of models with a strongly connected compartmental

graph is persistent for a wide class of rate functions. This result is based on the characterization of siphons in the Petri net of the system and the corresponding conserved quantities. The persistence property was then used in the stability analysis of compartmental models. It was shown that for strongly connected compartmental models, a unique equilibrium point exists within each stoichiometric compatibility class, and this equilibrium is asymptotically stable within each compatibility class even if the initial conditions are on the boundary of the nonnegative orthant (except for the two trivial boundary equlibria). Further work will be focused on the control related application of our results.

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