Normalish amenable subgroups of the R. Thompson groups

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Results in C^* algebras, of Matte Bon and Le Boudec, and of Haagerup and Olesen, apply to the R. Thompson groups $F \leq T \leq V$. These results together show that F is non-amenable if and only if T has a simple reduced C^* -algebra.

In further investigations into the structure of C^* -algebras, Breuillard, Kalantar, Kennedy, and Ozawa introduce the notion of a normalish subgroup of a group G. They show that if a group G admits no non-trivial finite normal subgroups and no *normalish* amenable subgroups then it has a simple reduced C^* -algebra. Our chief result concerns the R. Thompson groups F < T < V; we show that there is an elementary amenable group E < F [where here, $E \cong \ldots \wr \mathbb{Z} \wr \mathbb{Z} \wr \mathbb{Z}$] with E normalish in V.

The proof given uses a natural partial action of the group V on a regular language determined by a synchronising automaton in order to verify a certain stability condition: once again highlighting the existence of interesting intersections of the theory of V with various forms of formal language theory.

Keywords: Thompson's groups; amenable; C^* -simplicity; regular language; normalish.

1. Introduction

In this note we show that for the R. Thompson groups $F \leq T \leq V$ there is an elementary amenable group $E \leq F$ so that E is **normalish** in each of the groups F, T, and V.

This article is a fuller version of the note [6]: it has a large percentage of text in common, but it also gives some fuller proofs and discussion, and it corrects various minor typos appearing in [6].

1.1. General motivating background

Various weakenings of the notion of normal subgroup were introduced between 2014 and 2018 in order to obtain insight into the C^* -simplicity of the (reduced) group algebra $C_r^*(G)$ of a group G. This has had particular impact for infinite simple groups such as the R. Thompson groups T and V. The concept of a normalish subgroup of a group was introduced by in the seminal paper of Breuilliard, Kalantar, Kennedy, and Ozawa [8]. They show that a discrete group G with no non-trivial

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finite normal subgroups and no amenable normalish subgroups is C^* -simple. In that paper, they also obtain the just-previously-announced result of Haagerup and Olesen [11] that if the reduced group C^* -algebra $C^*_r(T)$ is simple, then F is non-amenable.

Meanwhile, Kennedy in [13] shows that a countable group G is C^* -simple (has simple reduced C^* -algebra) if and only if G admits no non-trivial amenable URS (uniformly recurrent subgroup). Using this, Le Boudec and Matte Bon in [15] show the converse of the stated Haagerup-Olesen result, if F is non-amenable, then the reduced C^* -algebra of T must be simple.

Indeed, for those interested in the question of the non-amenability of the R. Thompson group F, the focus has passed through the exploration of the uniformly recurrent subgroups of T to understanding the point stabilisers of the action of T on its Furstenberg boundary. Here, there are two possible cases, and F will be non-amenable precisely if these point stabilisers are trivial (see [15]). Despite this shift, we find the concept of normalish subgroups of "nearly simple" or simple groups like F and T to be of interest, and that is the focus of this note.

1.2. Core results

Let $G \leq H$ be groups. The group G is **normalish in** H if for any finite set of elements $\{c_1, c_2, \ldots, c_k\}$ the intersection

$$\bigcap_{i=1}^{\kappa} G^{c_i}$$

is infinite.

Our chief result is the following:

Theorem 1. There is an embedding of the elementary amenable group

$$^{\infty}(\mathbb{Z}\wr\mathbb{Z})=\ldots\wr\mathbb{Z})\wr\mathbb{Z}\wr\mathbb{Z}$$

into R. Thompson's group F so that the image group E is normalish in V.

Observe the corollary that E is then an amenable normalish subgroup of both F and of T as well.

1.3. Specific history of the core result

We should mention some other history related to this result. In [5] we showed the existence of an infinite direct sum of copies of \mathbb{Z} that could be found embedded as a normalish amenable subgroup of F, and discussed our conjecture (disproven here) that any normalish amenable subgroup of T should either contain an embedded subgroup isomorphic to R. Thompson's group F or to a non-abelian free subgroup. Meanwhile, the paper [15] shows that V contains an amenable normalish torsion group Λ : the subgroup of V consisting of those elements which are automorphisms of

the infinite rooted binary tree \mathcal{T}_2 . These automorphisms arise as finite compositions of the tree automorphisms that swap the two child vertices of any particular vertex (copying the dependent trees identically). The group Λ is normalish for reasons that are very similar to why our own group E is normalish, and it is a limit of finite groups hence elementary amenable. However, the group Λ is not a subgroup of Fnor of T.

1.4. An unexpected visitor? (Controlling synchronising automata)

A note on the proof: for experts on R. Thompson groups, the embedded copy of E that we find will clearly be normalish in V after short inspection. However, the technical proof of this requires a bit of work in that the conjugation action on our generators needs to not introduce too many breakpoints into our group elements, and also in that we need to have an infinite set \mathcal{X}_S of group elements in E so that \mathcal{X}_S is 'largely closed' under translations by arbitrary elements in V. The second task is the harder one if we are to avoid having subgroups isomorphic to R. Thompson's group F. We approach this by introducing a partial action of V on a regular language \mathcal{W} which is determined by a synchronising automaton, and we show that there is a sub-language S so that for any $v \in V$ and long enough word $w \in S$ (length depending on the choice of v), we have that w is in the image of the partial action of v on S. We then link this to a partial action of V on a corresponding infinite specified subset \mathcal{X}_S of E. It follows from the above considerations that the intersection of any finite collection of conjugates of the set \mathcal{X}_S (and therefore of E) remains an infinite set.

This is not the first time that connections between the theory of the group V and language theory have become apparent. Indeed, the definition of elements of V as acting on the finite prefixes of infinite binary strings leads one to consider connections between V and the theory of push-down automata, and this perspective is central to the proof of Lehnert and Schweitzer [16] that any finitely generated subgroup of V is a CoCF group (see [12]), and to the current form of Lehnert's conjecture, that this last statement is actually an "if and only if" (see [7]).

Thanks:

We would like to thank Adrienne Le Boudec for kind and informative conversations where he has helped the author of the present note to understand some of the amazing events that have transpired in the field of C^* -algebras over the last six years.

2. The interval and the circle as quotients of Cantor space, and some related language

Let $I := [0, 1] \subset \mathbb{R}$ represent the unit interval in the real numbers. Let $\mathfrak{C} := \{0, 1\}^{\omega}$ represent the Cantor space that arises as the infinite cartesian product of the discrete space $\{0, 1\}$ with itself, with the product indexed by the ordinal ω . As we will act

on our Cantor space from the right via prefix substitutions, we will express elements of Cantor space as left infinite strings, so a typical element \bar{x} of \mathfrak{C} will be written as $\bar{x} = \ldots x_2 x_1 x_0$ where each x_i is either a 0 or a 1. Note that in this usage, and for such left-infinite strings, we will refer to any finite rightmost contiguous substring as a **prefix** of the infinite string (and we will use the word prefix in this way as well when comparing finite strings, which we will formalise below). The monoid $\{0, 1\}^*$ of finite strings under the concatenation operator "^" (e.g., 00110^1001 = 001101001) will be central to our analysis and we might refer to an element of $\{0, 1\}^*$ as an **address**, for reasons which will become clear.

We give the monoid of finite words $\{0,1\}^*$ the **prefix-based partial ordering** as follows: if $p_1, p_2 \in \{0,1\}^*$ with $p_1 = x_j x_{j-1} \dots x_1 x_0$ and $p_2 = y_k y_{k-1} \dots y_1 y_0$ (where each x_i and y_i is in the set $\{0,1\}$ for each valid index i), we say $p_1 \leq p_2$ if and only if $j \leq k$ and for all indices $0 \leq i \leq j$ we have $x_i = y_i$. Recall that with this partial ordering, a **complete antichain** \mathcal{A} of $\{0,1\}^*$ is a finite set $\{p_1, p_2, \dots, p_k\}$ so that for each pair of distinct indices i and j we have that p_i and p_j are **incomparable** (written $p_1 \perp p_2$, and meaning that both $p_1 \not\leq p_2$ and $p_2 \not\leq p_1$ are true) and for any $w \in \{0,1\}^*$ we have some index r so that either $w \leq p_r$ or $p_r \leq w$.

The monoid $\{0,1\}^*$ with the partial order above can be naturally drawn as a rooted infinite binary tree, with its vertices being the elements of $\{0,1\}^*$, and where we draw an edge from a vertex r to a vertex s if $r \leq s$ and the length of s (denoted |s|) is one greater than the length of r. We will denote this tree as \mathcal{T}_2 and sketch a small neighbourhood of its root in the figure below (the tree \mathcal{T}_2 is often drawn so as to "open out" as one descends).

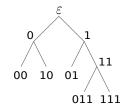


Fig. 1. A neighbourhood of the root ε of the tree \mathcal{T}_2

For any finite word $w = w_k w_{k-1} \dots w_1 w_0 \in \{0, 1\}^*$ we obtain the basic open set $\mathfrak{C}w$ for the topology of \mathfrak{C} . Specifically, $\mathfrak{C}w$ is the set of all points in Cantor space with prefix w:

$$\mathfrak{C}w = \{\bar{x}\,\hat{w} : \bar{x} \in \mathfrak{C}\}.$$

(Above, we are using the natural extension of the concatenation operator to allow us to concatenate a left infinite string with a finite string to its right.) We will refer to such basic open sets as cones, and for a given finite word $w \in \{0, 1\}^*$ the set $\mathfrak{C}w$ will be called the **cone at (address)** w. It is a standard fact that one can identify the Cantor space \mathfrak{C} with the boundary of \mathcal{T}_2 , or with the set of infinite descending paths in the tree (which correspond to infinite sequences of edge labels, if one labels each edge of \mathcal{T}_2 with a 0 or a 1, depending on the letter of the extension connecting the shorter address to the longer address.)

Recall there is a standard quotient map $q : \mathfrak{C} \twoheadrightarrow [0, 1]$, which we define fully here in order to give some practice with our right-to-left indexing notation. Let

$$\bar{x} = \dots x_2 x_1 x_0 \in \mathfrak{C}.$$

We have

$$(\bar{x})q := \sum_{i=0}^{\infty} x_i \cdot \frac{1}{2^{i+1}},$$

which we can think of as the ordinary map which interprets a real number in [0, 1] from its binary expansion (using the binary sequence moving leftward of the "binary point" as progressively smaller fractional parts of the unit).

We further recall that given any prefix $w = w_k w_{k-1} \dots w_1 w_0$ the map q identifies the two points $\overline{1}0w_k w_{k-1} \dots w_1 w_0$ and $\overline{0}1w_k w_{k-1} \dots w_1 w_0$. The resulting two-point equivalence classes map onto the dyadic rationals in $\mathbb{Z}[1/2] \cap (0,1) \subset \mathbb{R}$, and further, the cone $\mathfrak{C}w$ at w maps to the closed interval \mathcal{I}_w of radius $(1/2)^{k+2}$ centered at the diadic point d_w which is defined by the infinite sum

$$d_w := \sum_{i=0}^{\infty} w_i \cdot \frac{1}{2^{i+1}}.$$

where we set $w_{k+1} = 1$ and $w_m = 0$ for all m > k+1. For example, if w = 01, then k = 1 and we have $w_0 = 1$, $w_1 = 0$, $w_2 = 1$ and $w_m = 0$ for all m > 2. Then, d_{01} is computed as

$$d_{01} = \left(w_0 \cdot \frac{1}{2^1} + w_1 \cdot \frac{1}{2^2} + w_2 \cdot \frac{1}{2^3} + \mathbf{0}\right) = \left(1 \cdot \frac{1}{2} + 0 \cdot \frac{1}{4} + 1 \cdot \frac{1}{8}\right) = \frac{5}{8}$$

and the interval \mathcal{I}_{01} is of radius $1/(2^{1+2}) = 1/8$ centered at $d_{01} = 5/8$. In particular, we have $\mathcal{I}_{01} = [1/2, 3/4] = [5/8 - 1/8, 5/8 + 1/8]$.

For $w \in \{0,1\}^*$, we call the interval \mathcal{I}_w constructed as above the standard dyadic interval at address w (or "the standard dyadic interval centred at d_w "), noting that these intervals are naturally in a one-one correspondence with the words in the monoid $\{0,1\}^*$ (we set k = -1 when $w = \varepsilon$, the empty word, so that we produce the interval [0,1], that is, the closed interval of radius 1/2 centred at 1/2).

To obtain the circle as a quotient of Cantor space we add one further identification, that is, we identify the point $\dots 000 = \overline{0}0$ with the point $\dots 111 = \overline{1}1$, noting that this simply identifies the real numbers 0 and 1 from the interval I.

When working in the unit interval, we will mostly use the real number parameterisation of points, but sometimes it is convenient to name a point by one of its names arising from the map q^{-1} . Similarly, for points on the circle, we will use either the parameterisation arising from the quotient map $I \to I/(0 \sim 1) = \mathbb{R}/\mathbb{Z}$ (this is equivalent to applying the map $p: I \to \mathbb{S}^1$ given by $t \mapsto e^{2\pi i t}$ where we consider \mathbb{S}^1 as the unit circle in the complex plane) or, we will use the parameterisation arising from the map $q \cdot p: \mathfrak{C} \to \mathbb{S}^1$, where a point on the circle is referred to by one of its preimage left-infinite strings under the map $q \cdot p$.

Our group elements will act on the right, and induce permutations of the underlying sets of the spaces under consideration. We establish some notation for our context. Let Y be a set. We will use the notation Sym(Y) for the group of bijections from Y to itself. For any element $g \in Sym(Y)$ we define the **support of** g, written supt(g), as the set

$$\operatorname{supt}(q) := \{ y \in Y : yq \neq y \},\$$

that is, the set of points moved by g. In keeping with our right-actions notation, if $g, h \in \text{Sym}(Y)$, then the conjugate of g by h, denoted g^h , is the map $h^{-1}gh$. That is, we apply h^{-1} , then g, and finally h again. The following is a standard lemma from the theory of permutation groups.

Lemma 2. Let Y be a set, and $g, h \in Sym(Y)$. We have

$$\operatorname{supt}(g^h) = \operatorname{supt}(g)h.$$

It is quite useful in calculations to recall this way of thinking of the support of a conjugate: the support of g^h is the image of the support of g under the function h. As a fun aside for those readers not coming from a group theory background, this underlies one of the supporting principles for solving Rubik's Cube or other permutation puzzles: commutators have a fair chance of being elements of small support, and then conjugates of commutators are elements with targeted small and controlled impact.

3. The R. Thompson groups F < T < V

The Thompson groups F < T < V are groups of homeomorphisms which have been well studied. In this note, we generally take F, T, and V as each being groups of homeomorphisms of the Cantor space \mathfrak{C} .

3.1. Describing elements of F, T, and V

For two words $w_1, w_2 \in \{0, 1\}^*$ with $|w_1| > 0$ and $|w_2| > 0$ we define the **cone map** $\phi_{w_1,w_2} : \mathfrak{C}w_1 \to \mathfrak{C}w_2$ by the rule $\bar{x}w_1 \mapsto \bar{x}w_2$, for each point \bar{x} of \mathfrak{C} . It is immediate that this map is a homeomorphism from the Cantor space $\mathfrak{C}w_1$ to the Cantor space $\mathfrak{C}w_2$. Note that the map ϕ_{w_1,w_2} induces a map $\mathcal{I}_{w_1} \to \mathcal{I}_{w_2}$ which is a restriction of an affine map on the reals \mathbb{R} , and for this reason we might refer to ϕ_{w_1,w_2} as an "affine map" between the two subspaces of our larger Cantor space \mathfrak{C} . Note further that any such cone map ϕ_{w_1,w_2} is not just a homeomorphism from its domain to its

range but also that it has many extensions to homeomorphisms from $\mathfrak{C} \to \mathfrak{C}$, and we can think of ϕ_{w_1,w_2} as being a subset of a larger (if $w_1 \neq \varepsilon \neq w_2$) function from \mathfrak{C} to \mathfrak{C} (which we in turn consider as a subset of $\mathfrak{C} \times \mathfrak{C}$).

We are now in a position to define the R. Thompson groups F < T < V.

An element $g \in \text{Homeo}(\mathfrak{C})$ is an element of V if and only if we can write g as a **prefix replacement map**, as follows.

The element g is a prefix replacement map if and only if it admits some natural number n > 1, two complete antichains $D = \{a_1, a_2, \ldots, a_n\}$ and $R = \{r_1, r_2, \ldots, r_n\}$ for $\{0, 1\}^*$, and a bijection $\sigma : D \to R$, so that when restricted to any cone $\mathfrak{C}a_i$ (for valid index *i*), the map g restricts and co-restricts to the cone map $\phi_{a_i,a_i\cdot\sigma}$. In this context, we will write

$$g = (\{a_1, a_2, \dots, a_n\}, \{r_1, r_2, \dots, r_n\}, \sigma).$$

One could argue that the indices on the elements of D and R ought to be assigned so that the order preserving bijection determines g, so that σ becomes unnecessary. However, when composing many maps, it is often convenient to list the sets D and R in dictionary order, and then to determine σ as the requisite permutation. Thus, in typical usage the permutation σ is included explicitly. We observe in passing that any element of V admits infinitely many distinct prefix replacement maps representing it, but it is a standard exercise in R. Thompson theory that there is a unique, minimal (in the size of the set D) prefix-map representation of g.

We observe that any complete antichain $\{a_1, a_2, \ldots, a_n\}$ for $\{0, 1\}^*$ admits a natural left-to-right ordering \prec induced from the arrangement of the addresses a_i on the tree (this is simply the dictionary order, when we take $0 \prec 1$ and read our strings from right to left). An element $g \in V$ is in the subgroup F if and only if, when expressed as a prefix replacement map, the permutation σ preserves the ordering \prec . An element $g \in V$ is in the subgroup T if and only if, when expressed as a prefix replacement map, the permutation σ preserves the ordering \prec up to some cyclic rotation. The group V has no restriction on the type of permutations σ allowed. From this one can check that elements of F induce homeomorphisms of I through the quotient map q which are piecewise affine, respect the dyadic rationals, and where all slopes are powers of two and all breaks in slope occur over dyadic rationals. Similarly, the reader may check that elements of T induce homeomorphisms of \mathbb{S}^1 through the quotient $q \cdot p$ which are piecewise affine, respect the dyadic rationals, and where all slopes are powers of two and all breaks in slope occur over dyadic rationals.

A standard introductory reference for the general theory of the R. Thompson groups F, T, and V is the paper [9].

3.2. The element family X

We now single out a family

$$\mathcal{X} := \{x_w : w \in \{0, 1\}^*\}$$

of elements of V of specific interest to our discussion.

Given a word $w \in \{0, 1\}^*$, we specify the element x_w as the element of V which acts as the identity over the complement of the cone $\mathfrak{C}w$, and on the cone $\mathfrak{C}w$, acts according to the prefix map specified below (we only express the actual prefix substitutions here):

$$x_w := \begin{cases} 00^{\circ}w \mapsto 0^{\circ}w \\ 10^{\circ}w \mapsto 01^{\circ}u \\ 1^{\circ}w \mapsto 11^{\circ}w \end{cases}$$

In particular, the element x_w is the extension of the partial function

 $\phi_{00^{\widehat{}}w,0^{\widehat{}}w} \sqcup \phi_{10^{\widehat{}}w,01^{\widehat{}}w} \sqcup \phi_{1^{\widehat{}}w,11^{\widehat{}}w}$

by the identity map away from the cone $\mathfrak{C}w$.

Note that it is easy to extend the set $\{00^{\circ}w, 01^{\circ}w, 1^{\circ}w\}$ to a complete antichain $\{a_1, a_2, \ldots, a_{k-1}, 00^{\circ}w, 10^{\circ}w, 1^{\circ}w\}$ for $\{0, 1\}^*$ where |w| = k, and that in this case $\{a_1, a_2, \ldots, a_{k-1}, 0^{\circ}w, 01^{\circ}w, 11^{\circ}w\}$ is also a complete antichain for $\{0, 1\}^*$ (the set of addresses $\{a_i : 1 \leq i \leq k-1\}$ represents the minimal set of addresses one can use so that $\{a_i : i \in 1 \leq i \leq k-1\} \cup \{w\}$ is a complete antichain). Our map x_w acts as cone maps on each of the cones at the set of addresses $\{00^{\circ}w, 01^{\circ}w, 1^{\circ}w\}$, and otherwise takes each cone $\mathfrak{C}a_i$ to itself with the identity map.

It is easy to see that $\mathcal{X} \subset F$ so also $\mathcal{X} \subset T$ and $\mathcal{X} \subset V$. The figure below depicts the graphs of x_{ε} and x_{10} as homeomorphisms of I as examples.

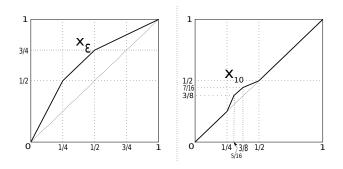


Fig. 2. The elements x_{ε} and x_{10} .

The family \mathcal{X} plays a central role in the dissertation [14] of Nayab Khalid, where Khalid works to provide a normal form for elements of F using \mathcal{X} as a natural infinite generating set. A key point is that a generator x_w has the effect of applying a rotation to the tree \mathcal{T}_2 at location w, and Khalid's normal forms potentially represent minimal length sequences of rotations (in the sense of Thurston, et al. in [17]) connecting one finite tree with k leaves to another. In theory, the algoritheorem in [14] runs in exponential time on the size of the tree, but, there are clear improvements that can be made and it remains unclear if the algoritheorem can be improved to run in polynomial time. Our small practice with it so far indicates that one can generally implement the algoritheorem by hand for reasonably small trees.

3.3. The element family \mathcal{G}

For each natural $n \in \mathbb{N}$, set $g_n := x_{(10)^n}$. The set $\mathcal{G} := \{g_i : i \in \mathbb{N}\}$ will be our second family of elements of F of interest. Note that we use \mathbb{N} to represent the natural numbers, which we take to be the non-negative integers.

We observe that Figure 2 also depicts g_0 and g_1 since $g_0 = x_{\varepsilon}$ and $g_1 = x_{10}$.

4. Realising $^{\infty}(\mathbb{Z} \wr \mathbb{Z})$ in F

Set $E := \langle \mathcal{G} \rangle \leq F \leq T \leq V$. For each index n, consider the group

$$W_n := \langle \{g_i : i \in \mathbb{N}, i < n\} \rangle$$

where for clarity we specify $W_0 = \{1_V\}$. It is immediate that $W_m \leq W_n$ when m < n. Direct calculation shows that $\sup(g_n) \cap \sup(g_n^{g_m}) = \emptyset$ whenever m < n (the content of the following lemma, which is not hard to prove). Specifically, following the arguments of [2,3] for natural index n we have $W_n \cong (\dots (\mathbb{Z} \setminus \mathbb{Z}) \setminus \mathbb{Z}) \cup \mathbb{Z}) \setminus \mathbb{Z}$ (with n appearances of \mathbb{Z} in this expression), which is a solvable group of derived length n. We immediately obtain $E \cong (\mathbb{Z} \setminus \mathbb{Z}) = \dots \in \mathbb{Z}) \setminus \mathbb{Z} \cup \mathbb{Z}$, as described in detail in [4]. As E admits a decomposition as a direct union of the solvable groups W_n , we obtain that E is elementary amenable (see Chou's paper [10] for details on the class of elementary amenable groups).

For what follows, set $G_{m,n} := \langle g_m, g_n \rangle$ for all natural numbers m < n.

Lemma 3. Let m < n be two natural numbers. We have

- (1) there is an isomorphism $G_{m,n} \cong G_{0,(n-m)}$ which is induced by a restriction map followed by a topological conjugacy,
- (2) $\operatorname{supt}(g_n) \cap \operatorname{supt}(g_n^{g_m}) = \emptyset$, and therefore
- (3) $G_{m,n} \cong \mathbb{Z} \wr \mathbb{Z}$.

Proof.

We first prove Point (1) that there is an isomorphism $G_{m,n} := \langle g_m, g_n \rangle \cong \langle g_0, g_{n-m} \rangle = G_{0,(n-m)}$ which is induced by a restriction map followed by a topological conjugacy.

To see this point, first observe that both of the elements g_n and g_m are supported wholly in the cone $\mathfrak{C}(10)^m$, so, the restriction of the maps g_n and g_m to the cone $\mathfrak{C}(10)^m$ results in an isomorphism of groups between the homeomorphism group $G_{m,n} = \langle g_m, g_n \rangle$, which acts on the Cantor space \mathfrak{C} , to a homeomorphism group $\widehat{G}_{m,n} = \langle \widehat{g}_m, \widehat{g}_n \rangle$ (these generators being the restrictions of the generators g_m and g_n respectively to the Cantor space $\mathfrak{C}(10)^m$), so that the group $\widehat{G}_{m,n}$ is a group of homeomorphisms of the Cantor space $\mathfrak{C}(10)^m$. Now, the homeomorphism $\theta_m : \mathfrak{C}(10)^m \to \mathfrak{C}$ which is induced by deleting the prefix $(10)^m$ from all points in the Cantor space $\mathfrak{C}(10)^m$ provides a topological conjugacy which induces an isomorphism from the group $\widehat{G}_{m,n}$ to the group $G_{0,n-m} = \langle g_0, g_{(n-m)} \rangle$, as under this topological conjugacy the image of the \widehat{g}_m is the element g_0 and the image of $\widehat{g}_{(n)}$ is the element $g_{(n-m)}$.

For Point (2), we observe that the restrictions applied in the argument for Point (1) only removed areas from the domain of the elements g_m and g_n where these elements already acted as the identity. Therefore the support of g_n and of $g_n^{g_m}$ will be disjoint if and only if the supports of the elements $g_{(n-m)}$ and of $g_{(n-m)}^{g_0}$ are disjoint. In particular, we have our result if we prove that for any positive integer k, we have $\sup(g_k) \cap \sup(g_k^{g_0}) = \emptyset$.

However, $g_0 = x_{\varepsilon}$, which acts over the cone $\mathfrak{C}10$ as a cone map, affinely taking the cone $\mathfrak{C}10$ rightward to the cone $\mathfrak{C}01$ by the prefix substitution $10 \mapsto 01$. Now, the support of g_k is contained in the cone $\mathfrak{C}(10)^k$, a subset of the cone $\mathfrak{C}(10)$. Direct calculation now shows that the cone $\mathfrak{C}(10)^k$ is carried affinely to the cone $\mathfrak{C}(10)^{k-1}(01)$ by g_0 , so Lemma 2 implies our result. We note in passing that we have shown that $g_k^{g_0} = x_{(10)^{k-1}01}$, or more specifically, that $x_{(01)^k}^{x_{\varepsilon}} = x_{(10)^{k-1}01}$, since our conjugator acted affinely.

For Point (3), recall Section 1.2.1 of [2], where an argument is given that two elements α_1 and α_2 of F generate a group isomorphic to $\mathbb{Z} \wr \mathbb{Z}$, with the element α_1 generating the top group of the wreath product. It happens that the element α_1 of that paper is the element we call x_{ε} here, while the element α_2 is the element we call x_{10} here. The proof of Section 1.2.1 essentially relies on only three facts: 1) the support of α_2 is contained in the support of α_1 , 2) every point in the support of α_1 is on an infinite orbit under the action of $\langle \alpha_1 \rangle$, and 3) the support of α_2 is moved entirely off of itself by α_1 . In our case with g_m and g_n , we again have these three conditions (with g_m playing the role of α_1), so we have our claimed Point (3). \Box

The discussion above indicates the following lemma.

Lemma 4. Let $v \in V$. There are $\{a_1, a_2, \ldots, a_n\}$ and $\{b_1, b_2, \ldots, b_n\}$, minimal cardinality (finite) antichains, together with a bijection σ between them, so that v can be described as the prefix replacement map

$$v = (\{a_1, a_2, \dots, a_n\}, \{b_1, b_2, \dots, b_n\}, \sigma).$$

If $u \in \{0,1\}^*$ then $x_{u^{\hat{-}}a_i}^v = x_{u^{\hat{-}}(a_i\sigma)}$ for all i.

Proof. By definition, we have $x_{u^{\uparrow}a_{i}}^{v} = v^{-1}x_{u^{\uparrow}a_{i}}v$. We note that the initial map v^{-1} restricts to a cone map from $\mathfrak{C}(a_{i}\sigma)$ to $\mathfrak{C}a_{i}$, that is, an affine map with image containing the support of $x_{u^{\uparrow}a_{i}}$. The action of $x_{u^{\uparrow}a_{i}}$ off the cone $\mathfrak{C}a_{i}$ is as the identity, and in general it acts as a prefix replacement map, which modifies only the prefixes which begin with $u^{\uparrow}a_{i}$ (so these modifications only appear in entries at indices larger than the length $|u^{\uparrow}a_{i}|$). Finally, v acts on the cone $\mathfrak{C}a_{i}$ by affinely returning it to $\mathfrak{C}(a_{i}\sigma)$ as a cone map, (it simply transforms the prefix a_{i} to the prefix $a_{i}\sigma$, and preserves all later entries (with index offset of size $|a_{i}\sigma| - |a_{i}|$) at larger indices,

for any point in the Cantor space $\mathfrak{C}(a_i)$). Therefore, $x_{u^{\hat{a}}a_i}^v = x_{u^{\hat{a}}(a_i\sigma)}$.

5. On partial actions

The proof of Lemma 4 suggests the well-known fact that V has a natural partial action on the addresses in $\{0, 1\}^*$. Let $v \in V$ and suppose that n is the minimal natural number so that there exist antichains $A = \{a_1, a_2, \ldots, a_n\}$ and $B = \{b_1, b_2, \ldots, b_n\}$ and a bijection σ between them so that v can be described as the prefix replacement map $v = (A, B, \sigma)$. The partial action of v on $\{0, 1\}^*$ is defined precisely on the set of words in $\{0, 1\}^*$ which admit one of the a_i as a prefix. Let us suppose $w \in \{0, 1\}^*$ and $w = u^{\hat{a}_i}$. We set $w \cdot v := u^{\hat{a}_i\sigma}$.

We can now re-express the result of Lemma 4 in terms of the partial action of V on the set $\{0, 1\}^*$.

Corollary 5. Let $v \in V$ and $w_1, w_2 \in \{0, 1\}^*$ so that $w_1 \cdot v = w_2$ under the partial action of V on $\{0, 1\}^*$. If $u \in \{0, 1\}^*$ then $x_{u \land w_1}^v = x_{u \land w_2}$.

That is, we see that the group V admits a partial action on the set \mathcal{X} which parallels its partial action on $\{0, 1\}^*$. We now work to understand the action of V on elements of the group E.

Our first step in understanding this partial action is to analyse a formal language.

5.1. A regular language and an action

Define the set $\mathcal{T} \subset \{0,1\}^*$ of **tokens** as follows:

$$\mathcal{T} := \{10^k, 01^k : k \in \mathbb{N}, k \neq 0\}.$$

We build a formal language \mathcal{W} over the alphabet $\{0, 1\}$ as follows. The language \mathcal{W} is the set of all words which decompose as $w = w_j \hat{w}_{j-1} \hat{\ldots} \hat{w}_1$ for some natural j, where each w_i is a token. The language \mathcal{W} is actually a regular language, which is recognised by the automaton \mathcal{A} depicted in Figure 3. The state q_0 of \mathcal{A} is both the start and accept state of \mathcal{A} .

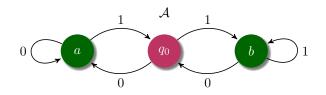


Fig. 3. The automaton \mathcal{A} which accepts the language \mathcal{W}

Recall that an automaton is a finite directed edge-labelled graph with a subset of its set of states called the start states of the automaton, and another subset of its

set of states called the accept states of the automaton. Then, the language accepted by the automaton is precisely the set of all finite-length words which arise as the concatenated edge-labels of some finite path in the automaton from a start state to an end state.

One can see that our formal language \mathcal{W} is indeed the language accepted by \mathcal{A} ; the paths which leave the state q_0 and then eventually return (exactly once) have labels of the form 10^k or 01^k , for some non-zero natural number k. In particular, the language accepted by \mathcal{A} is precisely the language of words built by concatenating tokens from \mathcal{T} .

We now set some terminology describing the structure of elements of the language \mathcal{W} . Observe firstly that the decomposition of any word in \mathcal{W} into a concatenation of tokens is unique. Therefore, for each $w \in \mathcal{W}$ we can define the **token length** of w as the number of tokens in its decomposition as a concatenation of tokens. Note that we index these tokens from right to left: $w = w_k \hat{w}_{k-1} \hat{\ldots} \hat{w}_2 \hat{w}_1$.

We now observe that the partial action of V on $\{0,1\}^*$ restricts to an action of $\langle g_0 \rangle$ on the set \mathcal{T} . Below, the proofs of Lemmas 6 and 8 follow by simple inductions on a basic calculation.

Lemma 6. For each integer k, the restriction of g_0^k to the cone $\mathfrak{C}10$ produces a cone map from the cone $\mathfrak{C}10$ to the cone $\mathfrak{C}w_k$, where w_k is given by the formula below:

$$w_k = 10 \cdot g_0^k = \begin{cases} 10^{|k-1|} & k \le 0\\ 01^k & k > 0. \end{cases}$$

Proof. This proof is a simple induction. Recall that $g_0 = x_{\varepsilon}$.

If k = 0 we observe that our formula works as $g_0^0 = 1_V$, which maps the cone $\mathfrak{C}10$ to the cone $\mathfrak{C}10$ by the identity map, which is a cone map. If k = 1 then g_0 takes $\mathfrak{C}10$ to $\mathfrak{C}01$ as a cone map, in accordance with the definiton of x_{ε} . For all $k \geq 2$, g_0^k acts as $x_{\varepsilon} \cdot x_{\varepsilon}^{k-1}$, so first as a cone map from $\mathfrak{C}10$ to $\mathfrak{C}01$, and then it will continue to act as x_{ε}^{k-1} on this resulting cone. However, the cone $\mathfrak{C}01$ is contained in the cone $\mathfrak{C}1$, and so the prefix replacement of x_{ε} here replaces the initial prefix 1 with the prefix 11, and this process repeats so inductively we have our desired result for all integers $k \geq 0$. For negative integers k, the argument follows as x_{ε}^{-1} replaces the prefix 0 with the prefix 00, and so inductively, the cone at 10 is carried by a cone map to the cone at $10^{|k-1|}$ by x_{ε}^k .

A translation of the above result is as follows.

Corollary 7. The partial action of V on the set $\{0,1\}^*$ restricts to a free, transitive action of $\langle g_0 \rangle$ on the set \mathcal{T} of tokens.

The following lemma simply extends the result of Lemma 6.

Lemma 8. For each integer k and natural number i, the restriction of g_i^k to the cone $\mathfrak{C}(10)^{i+1}$ produces a cone map from the cone $\mathfrak{C}(10)^{i+1}$ to the cone $\mathfrak{C}w_{i,k}$, where

 $w_{i,k}$ is given by the formula below:

$$w_{i,k} = (10)^{i+1} \cdot g_i^k = \begin{cases} 10^{|k-1|} (10)^i & k \le 0\\ 01^k (10)^i & k > 0. \end{cases}$$

Proof. The proof is similar to the proof of Lemma 6; $g_i = x_{(10)^i}$ acts as the identity off of the cone $\mathfrak{C}(10)^i$, and acts on the cone $\mathfrak{C}(10)^i$ in the same way that g_0 acts on the cone $\mathfrak{C}\varepsilon = \mathfrak{C}$ (this is essentially the content of the proof of Lemma 3(1). That is, the prefix $(10)^i$ is fixed by all powers of g_i , but the word $(10)^{i+1}$ is changed by g_i on the final token "10" (the "0" at index 2i and the "1" at index 2i + 1).

Lemma 8 has the following related corollary.

Corollary 9. Let *i* be a natural number. The partial action of *V* on the set $\{0,1\}^*$ restricts to a transitive and free action of $\langle g_i \rangle$ on the set of words $\{t^{(10)i} : t \in \mathcal{T}\}$.

6. Visiting the family \mathcal{X}

We now discuss the intersection of the group $E = \langle \mathcal{G} \rangle$ with the family \mathcal{X} .

Lemma 10. Let $w \in \mathcal{W}$, and $k \in \mathbb{N}$ so that w has token decomposition $w = w_k \, \hat{w}_{k-1} \, \dots \, \hat{w}_2 \, \hat{w}_1$. For each token w_i , let j_i be the integer so that $10 \cdot x_{\varepsilon}^{j_i} = w_i$ and also, recall that $x_{\varepsilon} = g_0$. If we set θ_w to be the product

$$\theta_w := g_{k-1}^{j_k} g_{k-2}^{j_{k-1}} \cdots g_1^{j_2} g_0^{j_1}$$

then we have

$$x_w = g_k^{\theta_w}$$

Proof. We construct θ_w as a product of (powers of) the generators in \mathcal{G} so as to create an element that modifies the prefix $(10)^k$ to the prefix w. To do this, we act on one token at a time, starting with the leftmost token (the k^{th} token w_k), and then working rightward to the first token w_1 . Progressively, each term in the product decomposition of θ_w acts on a cone containing the impact of the previous terms which have acted, the actions stack to create the following sequence of prefixes for the locations of the actions of the (partially) conjugated versions of $x_{(10)^k}$.

$$(10)^{k} \mapsto w_{k}(10)^{k-1} \mapsto w_{k}w_{k-1}(10)^{k-2} \mapsto \dots w_{k}w_{k-1}\dots w_{2}10 \mapsto w_{k}w_{k-1}\dots w_{2}w_{1}.$$

We therefore have the following corollary.

Corollary 11. Let $\mathcal{X}_{\mathcal{W}} := \{x_w : w \in \mathcal{W}\}$. Then $\mathcal{X}_{\mathcal{W}} \subset E$.

We now consider a special subset of \mathcal{W} . Set

$$\mathcal{S} := \{100^{\circ}w, 011^{\circ}w : w \in \{0, 1\}^*\}.$$

Lemma 12. The set S is a subset of W.

Proof. The automaton \mathcal{A} of Figure 3 has further properties not mentioned previously; it is highly connected and synchronising. These properties together mean that given any particular state (let us say q_0), there is a non-empty set of **synchronising words** W_{q_0} associated with q_0 so that, starting from any particular state s of the automaton and following a path labelled by any word in W_{q_0} , perforce, one will be lead to the state q_0 .

Note that the words 100 and 011 are synchronising words for the state q_0 ; no matter what state one starts in, after following the path labelled by the word 100 from that state, or the path labelled by the word 011 from that state, one arrives in the state q_0 (recall that we are reading these words from right-to-left!).

Thus, if we have some general word w and we append a suffix 011 or 100 to produce either $z = 011^{w}$ or $z = 100^{w}$ (that is, a general word z in S), then upon reading this resulting word on the automaton A starting from the start state q_0 , we will return to q_0 ; our word z is in the language W accepted by A.

In the proof above, we made essential use of the fact that our automaton was synchronising. The theory of synchronising automata has a long and interesting history, see [19, 18, 1], a woefully incomplete list of some of the many influential publications in this area.

Below, we will actually be interested in the subset of $\mathcal{X}_{\mathcal{W}}$ where the words involved come from \mathcal{S} . Set

$$\mathcal{X}_S := \{ x_w : w \in \mathcal{S} \}.$$

Corollary 13. The set \mathcal{X}_S is a subset of the group E.

We now consider how V interacts with the set \mathcal{X}_S under conjugation. The following lemma follows quickly by an application of Corollary 5.

Lemma 14. Let $v \in V$. There is a natural number n so that for all $w \in S$ with $|w| \ge n$ there is $z \in S$ so that $x_z^v = x_w$.

Proof. Let us assume that we can represent v by some prefix replacement map

$$v = (\{a_1, a_2, \dots, a_m\}, \{b_1, b_2, \dots, b_m\}, \sigma)$$

where we assume that n is at least three larger than the length of the largest string in the range antichain $\{b_1, b_2, \ldots, b_m\}$. We then set b as the prefix of w appearing in the set $\{b_1, b_2, \ldots, b_m\}$, and set $a = b\sigma^{-1} \in \{a_1, a_2, \ldots, a_m\}$. Then, $w = c^{b}$ where

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c is some string of length at least three, and we have $(c^{a}) \cdot v = (c^{b}) = w$, so that, in particular, if we take $z := c^{a}$, then Corollary 5 assures us that $x_{z}^{v} = x_{c^{b}} = x_{w}$.

But now, as the word c has length at least three, we see that it must end with the string 100 or the string 011 (since w has one of these two length three suffixes), and in particular, $z \in S$.

Thus, we have found that all sufficiently long strings w in S have that x_w is the conjugate image of x_z under v, for z another string in S.

Theorem 1: There is an embedding of the elementary amenable group

 $^{\infty}(\mathbb{Z}\wr\mathbb{Z})=\ldots\wr\mathbb{Z})\wr\mathbb{Z}\wr\mathbb{Z}$

into R. Thompson's group F so that the image group E is normalish in V.

Proof. The group E of this note is infinite and amenable. We can further see that for any finite set $C := \{v_1, v_2, \ldots, v_k\} \subset V$, the elements x_w , for $w \in S$ with w long enough (given by some particular integer dependent on the set C), all appear in all of the groups E^{v_i} . In particular, the intersection

$$\bigcap_{v_i \in C} E^i$$

is an infinite set, so that E is normalish in each of F, T, and V.

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