Algorithmic Graph Theory

A very brief and informal overview of some basic notations and definitions Martin Milanič, martin.milanic@upr.si

I. Graphs and graph parameters

- graph: G = (V, E) where V is a finite set of vertices and E is a subset of pairs of vertices (elements of E are referred to as edges)
 N(v) = {u ∈ V : uv ∈ E} is the set of neighbors of a vertex v, d(v) = d_G(v) = |N(v)| is its degree Δ(G): maximum degree of a vertex in G δ(G): minimum degree of a vertex in G
- K_n : a complete graph on *n* vertices; C_n : a cycle on *n* vertices; P_n : a path on *n* vertices; $K_{m,n}$: a complete bipartite graph with m + n vertices.
- independent set (stable set): a subset of pairwise non-adjacent vertices in a graph $\alpha(G)$ = independence number of *G*: maximum size of an independent set in *G*
- clique: a subset of pairwise adjacent vertices in a graph
 ω(G) = clique number of G: maximum size of a clique in G
- **dominating set**: a subset of vertices such that every vertex not in the set has a neighbor in the set

 $\gamma(G)$ = **domination number** of *G*: minimum size of a dominating set in *G*

- vertex cover: a subset of vertices such that every edge of the graph has at least one of its endpoints in the set
 τ(G) = vertex cover number of G: minimum size of a vertex cover in G
- matching: a subset of pairwise disjoint edges
 perfect matching: a matching covering all vertices of the graph
 ν(G) = matching number of G: maximum size of a matching in G
- *k*-(vertex) coloring: a mapping $c : V \to \{1, ..., k\}$ such that for all $u, v \in V$ such that $\{u, v\} \in E$, it holds $c(u) \neq c(v)$ *k*-colorable graph: a graph that admits a *k*-coloring $\chi(G)$ = chromatic number of G: minimum *k* such that G is *k*-colorable
- *k*-edge coloring: a mapping $c : E \to \{1, ..., k\}$ such that for all $e, f \in E$ such that $e \neq f$ and $e \cap f \neq \emptyset$, it holds that $c(e) \neq c(f)$ *k*-edge colorable graph: a graph that admits a *k*-edge coloring $\chi'(G)$ = chromatic index of *G*: minimum *k* such that *G* is *k*-edge colorable
- **list-chromatic number** of *G*: minimum *k* such that for all choices of sets S_v , $v \in V$, with $|S_v| \ge k$, there exists an *S*-coloring (a mapping $c : V \to \bigcup_{v \in V} S_v$ such that $c(v) \in S_v$ for all $v \in V$ and $c(u) \ne c(v)$ for all $uv \in E$)
- simplicial vertex: a vertex whose neighborhood is a clique

II. Graph operations

- H = (V', E') is a **subgraph** of G = (V, E) if $V' \subseteq V$ and $E' \subseteq E$
- H = (V', E') is an **induced subgraph** of G = (V, E) if $V' \subseteq V$ and $E' = \{e \in E \mid e \subseteq V'\}$; notation: H < G
- **disjoint union** of two graphs *G* and *H*: the graph obtained by adding to *G* a disjoint copy of *H* and no additional edges
- **join** of two graphs *G* and *H*: the graph obtained by adding to *G* a disjoint copy of *H* and all possible edges between *G* and *H*

- complement of a graph G = (V, E): a graph \overline{G} with vertex set V in which two vertices are adjacent if and only if they are non-adjacent in G
- the **line graph** of a graph *G* = (*V*,*E*) is the graph with vertex set *E* in which two distinct edges *e* and *f* are adjacent if and only if they have a common endpoint in *G*

III. Graph classes

- hereditary graph class: a set of graphs closed under vertex deletions (equivalently, closed under induced subgraphs)
- forest: a graph without cycles; tree: a connected forest
- **bipartite graph**: a graph such that there exists two disjoint sets *A* and *B* with $V = A \cup B$ such that $|e \cap A| = |e \cap B| = 1$ for every edge $e \in E$; equivalently, a 2-colorable graph
- **perfect graph**: a graph such that $\chi(H) = \omega(H)$ for all its induced subgraphs *H*; equivalently: a $\{C_5, C_7, \overline{C_7}, C_9, \overline{C_9}, \ldots\}$ -free graph
- threshold graph: a graph such that there exists non-negative weights $w : V \to \mathbb{R}_+$ and a threshold t such that for every $I \subseteq V$, $\sum_{v \in I} w(v) \leq t$ if and only if I is an independent set; equivalently: a $\{2K_2, C_4, C_5\}$ -free graph
- **split graph**: a graph that admits a partition of its vertex set into a clique and an independent set; equivalently: a $\{2K_2, C_4, C_5\}$ -free graph
- **cograph**: a graph that can be recursively built from copies of the one-vertex graph by iteratively applying the operations of disjoint union and join; equivalently: a *P*₄-free graph
- **chordal graph**: a graph in which every cycle of lenth at least 4 has a **chord** (an edge connecting two non-consecutive vertices of the cycle); equivalently: a {*C*₄, *C*₅,...}-free graph
- interval graph: intersection graph of closed intervals on the real line
- planar graph: a graph that can be drawn in the plane without edge crossings

IV. Graph problems

- INDEPENDENT SET: Given a graph *G* and an integer *k*, is $\alpha(G) \ge k$?
- CLIQUE: Given a graph *G* and an integer *k*, is $\omega(G) \ge k$?
- VERTEX COVER: Given a graph *G* and an integer *k*, is $\tau(G) \le k$?
- MATCHING: Given a graph *G* and an integer *k*, is $\nu(G) \ge k$?
- Dominating Set: Given a graph *G* and an integer *k*, is $\gamma(G) \leq k$?
- INDEPENDENT DOMINATING SET: Given a graph *G* and an integer *k*, is there an independent dominating set of size at most *k*?
- COLORABILITY: Given a graph *G* and an integer *k*, is $\chi(G) \leq k$?
- *k*-Colorability: Given a graph *G*, is $\chi(G) \leq k$?
- CHROMATIC INDEX: Given a graph *G* and an integer *k*, is $\chi'(G) \leq k$?
- Recognition of Graphs in *X*: Given a graph *G*, is $G \in X$?