# Cryptographically strong permutations from the butterfly structure 

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#### Abstract

Boomerang connectivity table is a new tool to characterize the vulnerability of cryptographic functions against boomerang attacks. Consequently, a cryptographic function is desired to have boomerang uniformity as low as its differential uniformity. Based on generalized butterfly structures recently introduced by Canteaut, Duval and Perrin, this paper presents infinite families of permutations of $\mathbb{F}_{2^{2 n}}$ for a positive odd integer $n$, which have the best known nonlinearity and boomerang uniformity 4 . Both open and closed butterfly structures are considered. The open butterflies, according to experimental results, appear not to produce permutations with boomerang uniformity 4 . On the other hand, from the closed butterflies we derive a condition on coefficients $\alpha, \beta \in \mathbb{F}_{2^{n}}$ such that the functions


$$
V_{i}(x, y):=\left(R_{i}(x, y), R_{i}(y, x)\right),
$$

where $R_{i}(x, y)=(x+\alpha y)^{2^{i}+1}+\beta y^{2^{i}+1}$ and $\operatorname{gcd}(i, n)=1$, permute $\mathbb{F}_{2^{n}}^{2}$ and have boomerang uniformity 4. In addition, experimental results for $n=3,5$ indicate that the proposed condition seems to cover all such permutations $V_{i}(x, y)$ with boomerang uniformity 4.

[^0]Keywords Permutations • Nonlinearity • Differential Uniformity • Boomerang Uniformity • Butterfly structure

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## 1 Introduction

Substitution boxes, known as S-boxes for short, are crucial nonlinear building blocks in modern block ciphers. In accordance with known attacks in the literature, S-boxes used in block ciphers are required to satisfy various cryptographic criteria, including high nonlinearity [6], low differential uniformity [17] and bijectivity. In Eurocrypt'18, Cid, Huang, Peyrin, Sasaki and Song [7] introduced a new tool of S-boxes, so-called the boomerang connectivity table (BCT), which analyzes the dependency between the upper part and lower part of a block cipher in a boomerang attack. This new tool quickly attracted researchers' interest in studying properties and bounds of BCT of cryptographic functions. Boura and Canteaut in [1] investigated the relation between entries in BCT and difference distribution table (DDT), and introduced the notion of the boomerang uniformity, which is the maximum value in BCT among all nonzero differences of inputs and outputs. They completely characterized the BCTs of 4 -bit S-boxes with differential uniformity 4 classified in [10], and also determined the boomerang uniformities of the inverse function and the Gold function. Later, $\mathrm{Li}, \mathrm{Qu}$, Sun and Li in [12] provided an equivalent formula to compute the boomerang uniformity of a cryptographic function. Using the new formula, they characterized the boomerang uniformity by means of the Walsh transform, and computed the boomerang uniformities of some permutations with low differential uniformity. Mesnager, Tang and Xiong considered the boomerang uniformity of quadratic permutations in [16], where they presented a characterization of quadratic permutations with boomerang uniformity 4 and showed that the boomerang uniformity of certain quadratic permutations is preserved under extended affine (EA) equivalence. Recently, Calderini and Villa [3] also investigated the boomerang uniformities of some non-quadratic permutations with differential uniformity 4. Very recently, Tian, Boura and Perrin [20] studied the boomerang uniformity of some popular constructions used for building large S-boxes, e.g. for eight variables from smaller ones.

It is shown that the boomerang uniformity of a cryptographic function is greater than or equal to its differential uniformity, and that the lowest possible boomerang uniformity 2 is achieved by almost perfect nonlinear (APN) functions [1, 7]. Clearly, APN permutations operating on even number of variables are most interesting. The problem of existence of APN permutations of $\mathbb{F}_{2^{2 n}}$ is referred to as the BIG APN problem in the community. Nonetheless, by far no other instance for this problem, except for the Dillon APN permutation of $\mathbb{F}_{2^{6}}$, has been found. Hence it is of great interest to construct permutations of $\mathbb{F}_{2^{2 n}}$ that have high nonlinearity, differential and boomerang uniformity 4. Up to now, there are only three infinite and inequivalent families of permutations over $\mathbb{F}_{2^{2 n}}$ that have boomerang uniformity 4 for odd integers $n \geq 1$ :
(1) $f(x)=x^{-1}$ over $\mathbb{F}_{2^{2 n}}[1]$;
(2) $f(x)=x^{2^{2 i}+1}$ over $\mathbb{F}_{2^{n}}$, where $\operatorname{gcd}(i, n)=1$ [1];
(3) $f(x)=\alpha x^{2^{2 s}+1}+\alpha^{2^{2 k}} x^{2^{-2 k}+2^{2 k+2 s}}$ over $\mathbb{F}_{2^{2 n}}$, where $n=3 k, 3 \nmid k, 3 \mid(k+s), \operatorname{gcd}(3 k, s)=1$, and $\alpha$ is a primitive element of $\mathbb{F}_{2^{2 n}}[16]$.

In Crypto'16, Perrin, Udovenko and Biryukov [19] investigated the only APN permutation over $\mathbb{F}_{2^{6}}$ [2] by means of reverse-engineering and proposed the open butterfly and the closed butterfly structures. A generalized butterfly structure was later proposed in [4]. The butterfly structures represent functions over $\mathbb{F}_{2^{n}}^{2}$ in terms of bivariate form. It is shown that the open butterfly structure produces permutations of $\mathbb{F}_{2^{n}}^{2}$, which are CCZ-equivalent to the functions that are derived from the closed structure and are in simpler forms [19]. Since differential uniformity is an invariant under CCZ-equivalence, one may consider to combine open and closed butterfly structures to construct permutations with low differential uniformity. As a matter of fact, by investigating differential uniformity of functions from the closed butterfly structure, researchers constructed several infinite families of differentially 4 -uniform permutations over $\mathbb{F}_{2^{n}}^{2}$ with the open butterfly structure $[4,5,8,14]$.

Motivated by recent works on the butterfly structure, this paper aims to construct infinite families of permutations with boomerang uniformity 4 from generalized butterfly structures. The main result of this paper is given as follows.

Theorem 1 Let $q=2^{n}$ with $n$ odd, $\operatorname{gcd}(i, n)=1$ and $R_{i}(x, y)=(x+\alpha y)^{2^{i}+1}+\beta y^{2^{i}+1}$ with $\alpha, \beta \in \mathbb{F}_{q}^{*}$, where $\mathbb{F}_{q}^{*}=\mathbb{F}_{q} \backslash\{0\}$. Then the function

$$
V_{i}(x, y):=\left(R_{i}(x, y), R_{i}(y, x)\right)
$$

from the closed butterfly structure permutes $\mathbb{F}_{q}^{2}$ and has boomerang uniformity 4 if $(\alpha, \beta)$ is taken from the following set

$$
\begin{equation*}
\Gamma=\left\{(\alpha, \beta) \in \mathbb{F}_{q}^{*} \times \mathbb{F}_{q}^{*} \mid \varphi_{2}^{2^{i}}=\varphi_{1} \varphi_{3}^{2^{i}-1} \text { and } \varphi_{3} \neq 0\right\} \tag{1}
\end{equation*}
$$

where $\varphi_{1}, \varphi_{2}, \varphi_{3}$ are given by

$$
\left\{\begin{array}{l}
\varphi_{1}=(\alpha+1)^{2^{i+1}+2}+\alpha^{2^{i}+2}+\alpha^{2^{i}}+\alpha \beta+\beta^{2}  \tag{2}\\
\varphi_{2}=(\alpha+1)^{2^{i+1}+2}+\alpha^{2^{i+1}+1}+\alpha+\alpha^{2^{i}} \beta+\beta^{2} \\
\varphi_{3}=(\alpha+1)^{2^{i+1}+2}+\beta^{2}
\end{array}\right.
$$

For the statement in Theorem 1, we will compute the boomerang uniformity by directly investigating the bivariate form $V_{i}(x, y)$, and prove the permutation property of $V_{i}(x, y)$ by examining its univariate polynomial representation over $\mathbb{F}_{q^{2}}$.

The rest of this paper is organized as follows. Section 2 firstly recalls the definitions of differential uniformity, boomerang uniformity, butterfly structure and introduces some auxiliary results. Sections 3 and 4 are devoted to proving the permutation property and the boomerang uniformity in Theorem 1, respectively. Finally, Section 5 draws a conclusion of our work.

## 2 Preliminaries

In this section, we assume $n$ is an arbitrary positive integer and $q=2^{n}$. Let $\operatorname{Tr}_{q}(\cdot)$ denote the absolute trace function over $\mathbb{F}_{q}$, i.e., $\operatorname{Tr}_{q}(x)=x+x^{2}+\cdots+x^{2^{n-1}}$ for any $x \in \mathbb{F}_{q}$. For any set $E$, the nonzero elements of $E$ is denoted by $E \backslash\{0\}$ or $E^{*}$.

### 2.1 Differential Uniformity and Boomerang Uniformity

The concept of differential uniformity was introduced to reveal the subtleties of differential attacks.

Definition 2 [17] Let $f(x)$ be a function from $\mathbb{F}_{q}$ to itself and $a, b \in \mathbb{F}_{q}$. The difference distribution table $(D D T)$ of $f(x)$ is given by a $q \times q$ table $D$, in which the entry for the $(a, b)$ position is given by

$$
D D T(a, b)=\#\left\{x \in \mathbb{F}_{q} \mid f(x+a)+f(x)=b\right\} .
$$

The differential uniformity of $f(x)$ is given by

$$
\Delta_{f}=\max _{a \in \mathbb{F}_{q}^{*}, b \in \mathbb{F}_{q}} D D T(a, b)
$$

It is straightforward for any function from $\mathbb{F}_{q}$ to itself, each entry in its DDT takes an even value and its differential uniformity is no less than 2. A function with the minimum possible differential uniformity 2 is called an almost perfect nonlinear (APN) function.

The concept of boomerang connectivity table of a permutation $f$ from $\mathbb{F}_{2}^{n}$ to itself was introduced in [7], which clearly is also suitable for the case $\mathbb{F}_{2^{n}}$. Later, Boura and Canteaut introduced the concept of the boomerang uniformity, which is defined by the maximum value in BCT excluding the first row and column.

Definition 3 [1, 7] Let $f$ be an invertible function from $\mathbb{F}_{q}$ to itself and $a, b \in \mathbb{F}_{q}$. The boomerang connectivity table (BCT) of $f$ is given by a $q \times q$ table, in which the entry for the $(a, b)$ position is given by

$$
\begin{equation*}
B C T(a, b)=\#\left\{x \in \mathbb{F}_{q}: f^{-1}(f(x)+b)+f^{-1}(f(x+a)+b)=a\right\} \tag{3}
\end{equation*}
$$

The boomerang uniformity of $f$ is defined by

$$
\delta_{f}=\max _{a, b \in \mathbb{F}_{q}^{*}} B C T(a, b)
$$

It is shown in $[1,7]$ that $B C T(a, b) \geq D D T(a, b)$ for any $a, b$ in $\mathbb{F}_{q}$. In [12], Li, Qu , Sun and Li presented an equivalent formula to compute BCT and the boomerang uniformity without knowing $f^{-1}(x)$ and $f(x)$ simultaneously as follows.

Proposition 4 [12] Let $q=2^{n}$ and $f(x) \in \mathbb{F}_{q}[x]$ be a permutation polynomial over $\mathbb{F}_{q}$. Then the $B C T$ of $f(x)$ can be given by a $q \times q$ table BCT, in which the entry $B C T(a, b)$ for the $(a, b)$ position is given by the number of solutions $(x, y)$ in $\mathbb{F}_{q} \times \mathbb{F}_{q}$ of the following equation system.

$$
\left\{\begin{array}{l}
f(x+a)+f(y+a)=b  \tag{4}\\
f(x)+f(y)=b
\end{array}\right.
$$

Equivalently, the boomerang uniformity of $f(x)$, given by $\delta_{f}$, is the maximum number of solutions in $\mathbb{F}_{q} \times \mathbb{F}_{q}$ of (4) as a, b run through $\mathbb{F}_{q}^{*}$.

Let $f$ be a quadratic function from $\mathbb{F}_{q}$ to itself with $f(0)=0$. The associated symmetric bilinear mapping is given by $S_{f}(x, y)=f(x+y)+f(x)+f(y)$, where $x, y \in \mathbb{F}_{q}$. For any $a \in \mathbb{F}_{q}$, define

$$
\operatorname{Im}_{f, a}=\left\{S_{f}(a, x): x \in \mathbb{F}_{q}\right\}
$$

Very recently, Mesnager, Tang and Xiong [16] presented a characterization about quadratic permutations with boomerang uniformity 4 using the new formula (4).

Lemma 5 [16] Let $q=2^{n}$ and $f$ be a quadratic permutation of $\mathbb{F}_{q}$ with differential uniformity 4. Then the boomerang uniformity of $f$ equals 4 if and only if $\operatorname{Im}_{f, a}=\operatorname{Im}_{f, b}$ for any $a, b \in \mathbb{F}_{q}^{*}$ satisfying $S_{f}(a, b)=0$.

### 2.2 The Butterfly Structure

In Crypto'16, Perrin, Udovenko and Biryukov [19] analyzed the only known APN permutation over $\mathbb{F}_{2^{6}}$ [2] and discovered that the APN permutation over $\mathbb{F}_{2^{6}}$ has a simple decomposition relying on $x^{3}$ over $\mathbb{F}_{2^{3}}$. Based on the power permutation $x^{e}$ over $\mathbb{F}_{2^{n}}$, they presented the open butterfly structure and the closed butterfly structure, which were later generalized in [4].

Definition 6 [19] Let $q=2^{n}$ and $\alpha \in \mathbb{F}_{q}$, e be an integer such that $x^{e}$ is a permutation over $\mathbb{F}_{q}$ and $R_{k}[e, \alpha]$ be the keyed permutation

$$
R_{k}[e, \alpha](x)=(x+\alpha k)^{e}+k^{e}
$$

The following functions

$$
\begin{gathered}
H_{e}^{\alpha}(x, y)=\left(R_{R_{y}[e, \alpha](x)}^{-1}(y), R_{y}[e, \alpha](x)\right), \\
V_{e}^{\alpha}(x, y)=\left(R_{y}[e, \alpha](x), R_{x}[e, \alpha](y)\right)
\end{gathered}
$$

are called the open butterfly structure and closed butterfly structure respectively.
Definition 7 [4] Let $q=2^{n}$ and $R(x, y)$ be a bivariate polynomial of $\mathbb{F}_{q}$ such that $R_{y}: x \rightarrow R(x, y)$ is a permutation of $\mathbb{F}_{q}$ for all $y$ in $\mathbb{F}_{q}$. The closed butterfly $V_{R}$ is the function of $\mathbb{F}_{q}^{2}$ defined by

$$
V_{R}(x, y)=(R(x, y), R(y, x))
$$

and the open butterfly $H_{R}$ is the permutation of $\mathbb{F}_{q}^{2}$ defined by

$$
H_{R}(x, y)=\left(R_{R_{y}^{-1}(x)}(y), R_{y}^{-1}(x)\right)
$$

where $R_{y}(x)=R(x, y)$ and $R_{y}^{-1}\left(R_{y}(x)\right)=x$ for any $x, y$.
Define a bivariate polynomial

$$
R_{i}(x, y)=(x+\alpha y)^{2^{i}+1}+\beta y^{2^{i}+1}, \quad \operatorname{gcd}(i, n)=1, \alpha, \beta \in \mathbb{F}_{q}
$$

Since $n$ is odd, it is clear that the mapping $x \mapsto R_{i}(x, y)$ is a permutation of $\mathbb{F}_{q}$ for any fixed $y \in$ $\mathbb{F}_{q}$. According to experimental results, the permutation $H_{R_{i}}(x, y)$ from $R_{i}(x, y)$ and the open butterfly
structure seems not to have boomerang uniformity 4 of $\mathbb{F}_{2^{3}}^{2}$. Hence this paper concentrates on the closed butterfly structure.

Lemma 8 [14] Let $n$ be odd, $q=2^{n}$, $i$ be an integer with $\operatorname{gcd}(i, n)=1, \alpha, \beta \in \mathbb{F}_{q}^{*}$ and $\beta \neq(\alpha+1)^{2^{i}+1}$. Then the function

$$
V_{i}:=\left(R_{i}(x, y), R_{i}(y, x)\right) \text { with } R_{i}(x, y)=(x+\alpha y)^{2^{i}+1}+\beta y^{2^{i}+1}
$$

has differential uniformity at most 4.
Recall that the boomerang uniformity of a function is no less than its differential uniformity. The result about $V_{i}(x, y)$ in Lemma 8 motivates our study on the coefficients $\alpha, \beta$ in $\mathbb{F}_{q}^{*}$ that can further result in permutations $V_{i}(x, y)$ with boomerang uniformity 4.

### 2.3 Useful Lemmas

This subsection summarizes some lemmas that will be used for proving the permutation property of the function in Theorem 1.

Lemma 9 [18, 21, 22] Pick $d, r>0$ with $d \mid(q-1)$, and let $h(x) \in \mathbb{F}_{q}[x]$. Then $f(x)=x^{r} h\left(x^{(q-1) / d}\right)$ permutes $\mathbb{F}_{q}$ if and only if both
(1) $\operatorname{gcd}(r,(q-1) / d)=1$ and
(2) $g(x)=x^{r} h(x)^{(q-1) / d}$ permutes $\mu_{d}$, where $\mu_{d}:=\left\{x \in \mathbb{F}_{q}: x^{d}=1\right\}$.

Let the unit circle of $\mathbb{F}_{q^{2}}$ be defined by

$$
\mu_{q+1}:=\left\{x \in \mathbb{F}_{q^{2}}: x^{q+1}=1\right\}
$$

The unit circle of $\mathbb{F}_{q^{2}}$ has the following relation with the finite field $\mathbb{F}_{q}$.
Lemma 10 [9] Let $\gamma$ be any fixed element in $\mathbb{F}_{q^{2}} \backslash \mathbb{F}_{q}$. Then we have

$$
\mu_{q+1} \backslash\{1\}=\left\{\frac{x+\gamma}{x+\gamma^{q}}: x \in \mathbb{F}_{q}\right\}
$$

The following lemma is about the solutions of a linear equation. The proof is easy and we omit it.
Lemma 11 Let $q=2^{n}$ and $\operatorname{gcd}(i, n)=1$. Then for any $a \in \mathbb{F}_{q}$, the equation $x^{2^{i}}+x=a$ has solutions in $\mathbb{F}_{q}$ if and only if $\operatorname{Tr}_{q}(a)=0$. Moreover, when $\operatorname{Tr}_{q}(a)=0$, the equation $x^{2^{i}}+x=a$ has exactly two solutions $x=x_{0}, x_{0}+1$ in $\mathbb{F}_{q}$.

Lemma 12 [15] Let $\mathbb{R}$ be a commutative ring with identity. The Dickson polynomial $D_{k}(x, a)$ of the first kind of degree $k$

$$
D_{k}(x, a)=\sum_{j=0}^{\left\lfloor\frac{k}{2}\right\rfloor} \frac{k}{k-j}\binom{k-j}{j}(-a)^{j} x^{k-2 j}
$$

has the following properties:
(1) $D_{k}\left(x_{1}+x_{2}, x_{1} x_{2}\right)=x_{1}^{k}+x_{2}^{k}$, where $x_{1}, x_{2}$ are two indeterminates;
(2) $D_{k+2}(x, a)=x D_{k+1}(x, a)-a D_{k}(x, a)$;
(3) $D_{k \ell}(x, a)=D_{k}\left(D_{\ell}(x, a), a^{\ell}\right)$;
(4) if $\mathbb{R}=\mathbb{F}_{2^{n}}$, then $D_{2^{i}}(x, a)=x^{2^{i}}$.

By the above lemma, the Dickson polynomial of degree $k=2^{i}-1$ over $\mathbb{F}_{2^{n}}$ can be explicitly given.

Lemma 13 For any positive integer $i$ and element $a \in \mathbb{F}_{2^{n}}$,

$$
\begin{equation*}
D_{2^{i}-1}(x, a)=\sum_{j=0}^{i-1} a^{2^{j}-1} x^{2^{i}-2^{j+1}+1} \tag{5}
\end{equation*}
$$

Proof We prove the statement by induction. It is clear that (5) holds for $i=1$ since $D_{1}(x, a)=x$. Suppose that (5) holds for $i-1$, namely,

$$
\begin{equation*}
D_{2^{i-1}-1}(x, a)=\sum_{j=0}^{i-2} a^{2^{j}-1} x^{2^{i-1}-2^{j+1}+1} \tag{6}
\end{equation*}
$$

By Lemma 12 (3) and (4), we have

$$
\begin{aligned}
D_{2^{i}-2}(x, a) & =D_{2^{i-1}-1}\left(D_{2}(x, a), a^{2}\right) \\
& =D_{2^{i-1}-1}\left(x^{2}, a^{2}\right) \\
& =\sum_{j=0}^{i-2} a^{2\left(2^{j}-1\right)} x^{2\left(2^{i-1}-2^{j+1}+1\right)} \\
& =\sum_{j=1}^{i-1} a^{2^{j}-2} x^{2^{i}-2^{j+1}+2} .
\end{aligned}
$$

In addition, according to Lemma 12 (2),

$$
D_{2^{i}}(x, a)=x D_{2^{i}-1}(x, a)+a D_{2^{i}-2}(x, a) .
$$

Thus,

$$
\begin{aligned}
D_{2^{i}-1}(x, a) & =x^{-1}\left(x^{2^{i}}+a D_{2^{i}-2}(x, a)\right) \\
& =x^{-1}\left(x^{2^{i}}+a \sum_{j=1}^{i-1} a^{2^{j}-2} x^{2^{i}-2^{j+1}+2}\right) \\
& =\sum_{j=0}^{i-1} a^{2^{j}-1} x^{2^{i}-2^{j+1}+1}
\end{aligned}
$$

which implies that (5) holds for the $i$ case. Therefore, the desired conclusion follows.

Let $\gamma$ be a primitive element of $\mathbb{F}_{2^{2}}$, i.e, $\gamma^{2}=\gamma+1$. Let $n$ be a positive odd integer and $q=2^{n}$. The finite field $\mathbb{F}_{q^{2}}=\mathbb{F}_{q}(\gamma)$ and the basis $1, \gamma$ of $\mathbb{F}_{q^{2}}$ over $\mathbb{F}_{q}$ induces a one-to-one correspondence between $\mathbb{F}_{q}^{2}$
and $\mathbb{F}_{q^{2}}$ as follows:

$$
z=x+\gamma y \leftrightarrow(x, y)=\left(\gamma^{2} z+\gamma z^{q}, z^{q}+z\right) .
$$

According to the one-to-one correspondence between $\mathbb{F}_{q}^{2}$ and $\mathbb{F}_{q^{2}}$, the closed butterfly structure $V_{i}(x, y)=$ $\left(R_{i}(x, y), R_{i}(y, x)\right)$ over $\mathbb{F}_{q}^{2}$ can be expressed in a univariate $z=x+\gamma y$ as

$$
V_{i}(z):=R_{i}(x, y)+\gamma R_{i}(y, x) \text { with } x=\gamma^{2} z+\gamma z^{q}, y=z^{q}+z .
$$

By substituting $z$ with $\gamma z$ when $i$ is odd (resp. $\gamma^{2} z$ when $i$ is even), the above univariate polynomial can be transformed into

$$
\begin{equation*}
f_{i}(z)=\epsilon_{1} z^{q \cdot\left(2^{i}+1\right)}+\epsilon_{2} z^{q \cdot 2^{i}+1}+\epsilon_{3} z^{z^{i}+q}+\epsilon_{4} z^{2^{i}+1} \tag{7}
\end{equation*}
$$

where the coefficients

$$
\left(\epsilon_{1}, \epsilon_{2}, \epsilon_{3}, \epsilon_{4}\right)= \begin{cases}\left(\varepsilon_{1}, \varepsilon_{2}, \varepsilon_{3}, \varepsilon_{4}\right), & \text { for even } i  \tag{8}\\ \left(\varepsilon_{3}, \varepsilon_{4}, \varepsilon_{1}, \varepsilon_{2}\right), & \text { for odd } i\end{cases}
$$

with

$$
\left\{\begin{array}{l}
\varepsilon_{1}=\alpha^{2^{i}}+\alpha+1  \tag{9}\\
\varepsilon_{2}=\alpha^{2^{i}+1}+\alpha+\beta+1 \\
\varepsilon_{3}=\alpha^{2^{i}+1}+\alpha^{2^{i}}+\beta+1 \\
\varepsilon_{4}=\alpha^{2^{i}+1}+\alpha^{2^{i}}+\alpha+\beta
\end{array}\right.
$$

Further, we define

$$
\left\{\begin{array}{l}
\varphi_{1}=\epsilon_{1} \epsilon_{3}+\epsilon_{2} \epsilon_{4}  \tag{10}\\
\varphi_{2}=\epsilon_{1} \epsilon_{2}+\epsilon_{3} \epsilon_{4} \\
\varphi_{3}=\epsilon_{1}^{2}+\epsilon_{2}^{2}+\epsilon_{3}^{2}+\epsilon_{4}^{2} \\
\varphi_{4}=\epsilon_{1}^{2}+\epsilon_{4}^{2}
\end{array}\right.
$$

It's easy to check that $\varphi_{i}$ 's, $i=1,2,3$, match the ones defined in (2).
At the end of this section, we provide a lemma about some properties of the elements $\varphi_{i}$ 's which are characterized in Theorem 1. This result will be heavily used in the proof of the main theorem.

Lemma 14 Let $q=2^{n}$ with $n$ odd and $\operatorname{gcd}(i, n)=1$. Let $\varphi_{1}, \varphi_{2}, \varphi_{3}, \varphi_{4}$ be defined by (10) satisfying $\varphi_{2}^{2^{i}}=\varphi_{1} \varphi_{3}^{2^{i}-1}$ and $\varphi_{3} \neq 0$. For $\alpha, \beta \in \mathbb{F}_{q}^{*}$, they have the following properties:
(1) $\left(\varphi_{1}+\varphi_{3}\right)\left(\varphi_{2}+\varphi_{3}\right)\left(\varphi_{3}+\varphi_{4}\right) \varphi_{4} \neq 0$ and $\left(\frac{\varphi_{3}}{\varphi_{2}+\varphi_{3}}\right)^{2^{i}}=\frac{\varphi_{3}}{\varphi_{1}+\varphi_{3}}$;
(2) when $i$ is even, $\operatorname{Tr}_{q}\left(\frac{\varphi_{4}}{\varphi_{3}}\right)=1$; moreover, the equation

$$
x^{2^{i}}+x+\frac{\varphi_{3}+\varphi_{4}}{\varphi_{3}}=0
$$

has two solutions $\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha$ and $\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha+1$ in $\mathbb{F}_{q} ;$
(3) when $i$ is odd, $\operatorname{Tr}_{q}\left(\frac{\varphi_{4}}{\varphi_{3}}\right)=0$;
(4) $\operatorname{Tr}_{q}\left(\frac{\varphi_{2}}{\varphi_{3}}\right)=0$.

Proof Since

$$
\left\{\begin{array}{l}
\varphi_{1}=\alpha^{2^{i+1}+2}+\alpha^{2^{i+1}}+\alpha^{2^{i}+2}+\alpha^{2^{i}}+\alpha^{2}+\alpha \beta+\beta^{2}+1  \tag{11}\\
\varphi_{2}=\alpha^{2^{i+1}+2}+\alpha^{2^{i+1}+1}+\alpha^{2^{i+1}}+\alpha^{2}+\alpha+\alpha^{2^{i}} \beta+\beta^{2}+1 \\
\varphi_{3}=\alpha^{2^{i+1}+2}+\alpha^{2^{i+1}}+\alpha^{2}+\beta^{2}+1
\end{array}\right.
$$

and

$$
\varphi_{4}= \begin{cases}\alpha^{2^{i+1}+2}+\beta^{2}+1 & \text { for even } i  \tag{12}\\ \alpha^{2^{i+1}}+\alpha^{2} & \text { for odd } i\end{cases}
$$

it is clear that

$$
\begin{equation*}
\varphi_{1}+\varphi_{2}=\alpha^{2^{i}+2}+\alpha^{2^{i}}+\alpha \beta+\alpha^{2^{i+1}+1}+\alpha^{2^{i}} \beta+\alpha=\left(\alpha^{2^{i}}+\alpha\right)\left(\alpha^{2^{i}+1}+\beta+1\right) \tag{13}
\end{equation*}
$$

and

$$
\begin{equation*}
\left\{\varphi_{4}, \varphi_{3}+\varphi_{4}\right\}=\left\{\left(\alpha^{2^{i}}+\alpha\right)^{2},\left(\alpha^{2^{i}+1}+\beta+1\right)^{2}\right\} \tag{14}
\end{equation*}
$$

(1) It follows from (13) and (14) that

$$
\varphi_{4}\left(\varphi_{3}+\varphi_{4}\right)=\left(\varphi_{1}+\varphi_{2}\right)^{2}=\left(\alpha^{2^{i}}+\alpha\right)^{2}\left(\alpha^{2^{i}+1}+\beta+1\right)^{2}
$$

By the equality $\varphi_{2}^{2^{i}}=\varphi_{1} \varphi_{3}^{2^{i}-1}$, if either $\varphi_{4}\left(\varphi_{3}+\varphi_{4}\right)=0$ or $\left(\varphi_{1}+\varphi_{3}\right)\left(\varphi_{2}+\varphi_{3}\right)=0$, then $\varphi_{1}+\varphi_{2}=0$ and $\varphi_{1}=\varphi_{2}=\varphi_{3}$. The equation $\varphi_{1}+\varphi_{2}=\left(\alpha^{2^{i}}+\alpha\right)\left(\alpha^{2^{i}+1}+\beta+1\right)=0$ implies $\beta=\alpha^{2^{i}+1}+1$ or $\alpha^{2^{i}}+\alpha=0$. In fact, if $\beta=\alpha^{2^{i}+1}+1$, then $\varphi_{1}+\varphi_{3}=\alpha^{2^{i}+2}+\alpha^{2^{i}}+\alpha \beta=\alpha^{2^{i}}+\alpha=0$. Thus we always have $\alpha^{2^{i}}+\alpha=0$, equivalently $\alpha=0,1$. This implies $\varphi_{1}+\varphi_{3}=\alpha^{2^{i}+2}+\alpha^{2^{i}}+\alpha \beta=\alpha \beta=0$, which is in contradiction with the assumption $\alpha \beta \neq 0$.

In addition, it is clear that the equality $\varphi_{2}^{2^{i}}=\varphi_{1} \varphi_{3}^{2^{i}-1}$ implies

$$
\left(\frac{\varphi_{3}}{\varphi_{2}+\varphi_{3}}\right)^{2^{i}}=\frac{\varphi_{3}}{\varphi_{1}+\varphi_{3}}
$$

(2) From (11) and (12), we have

$$
\left\{\begin{array}{l}
\varphi_{1}+\varphi_{3}=\alpha^{2^{i}+2}+\alpha^{2^{i}}+\alpha \beta  \tag{15.1}\\
\varphi_{2}+\varphi_{3}=\alpha^{2^{i+1}+1}+\alpha^{2^{i}} \beta+\alpha \\
\varphi_{3}+\varphi_{4}=\alpha^{2^{i+1}}+\alpha^{2}
\end{array}\right.
$$

It is easy to verify that

$$
\begin{equation*}
\alpha\left(\varphi_{2}+\varphi_{3}\right)+\alpha^{2^{i}}\left(\varphi_{1}+\varphi_{3}\right)=\varphi_{3}+\varphi_{4} . \tag{16}
\end{equation*}
$$

Moreover, using $\left(\frac{\varphi_{3}}{\varphi_{2}+\varphi_{3}}\right)^{2^{i}}=\frac{\varphi_{3}}{\varphi_{1}+\varphi_{3}}$, we have

$$
\begin{equation*}
\frac{\varphi_{3}+\varphi_{4}}{\varphi_{3}}=\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha+\frac{\varphi_{1}+\varphi_{3}}{\varphi_{3}} \alpha^{2^{i}}=\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha+\left(\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha\right)^{2^{i}} \tag{17}
\end{equation*}
$$

Thus,

$$
\operatorname{Tr}_{q}\left(\frac{\varphi_{3}+\varphi_{4}}{\varphi_{3}}\right)=\operatorname{Tr}_{q}\left(\frac{\varphi_{4}}{\varphi_{3}}\right)+\operatorname{Tr}_{q}(1)=0
$$

Furthermore, from Lemma 11, the solutions in $\mathbb{F}_{q}$ of $x^{2^{i}}+x=\frac{\varphi_{3}+\varphi_{4}}{\varphi_{3}}$ are $\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha$ and $\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha+1$.
(3) From the expressions of $\varphi_{3}, \varphi_{4}$ in (11), (12), it is clear that the values of $\varphi_{4}$ for even $i$ and odd $i$ add up to $\varphi_{3}$. The fact that $\operatorname{Tr}_{q}(1)=1$ for odd integer $n$ implies that the values of $\operatorname{Tr}_{q}\left(\frac{\varphi_{4}}{\varphi_{3}}\right)$ for even $i$ and odd $i$ add up to 1 . The desired assertion directly follows from (2) of this lemma.
(4) From (13) and (14), it is easily seen that

$$
\begin{equation*}
\left(\frac{\varphi_{1}}{\varphi_{3}}\right)^{2}+\left(\frac{\varphi_{2}}{\varphi_{3}}\right)^{2}=\left(\frac{\varphi_{4}}{\varphi_{3}}\right)^{2}+\frac{\varphi_{4}}{\varphi_{3}} . \tag{18}
\end{equation*}
$$

Plugging $\frac{\varphi_{1}}{\varphi_{3}}=\left(\frac{\varphi_{2}}{\varphi_{3}}\right)^{2^{i}}$ into Eq. (18), we get

$$
\left(\frac{\varphi_{2}}{\varphi_{3}}\right)^{2^{i+1}}+\left(\frac{\varphi_{2}}{\varphi_{3}}\right)^{2}=\left(\frac{\varphi_{4}}{\varphi_{3}}\right)^{2}+\frac{\varphi_{4}}{\varphi_{3}}=\left(\frac{\varphi_{3}+\varphi_{4}}{\varphi_{3}}\right)^{2}+\frac{\varphi_{3}+\varphi_{4}}{\varphi_{3}}
$$

By the relation between $\varphi_{3}$ and $\varphi_{4}$ for even and odd $i$, it is clear that the expression on the right side of the above equation is independent of the parity of the integer $i$. W.L.O.G., we can assume that $i$ is even, since the case $i$ odd can be proved by just replacing $\varphi_{4}$ by $\varphi_{3}+\varphi_{4}$. Together with (17), we have

$$
\left(\frac{\varphi_{2}}{\varphi_{3}}\right)^{2^{i+1}}+\left(\frac{\varphi_{2}}{\varphi_{3}}\right)^{2}=\left(\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha+\left(\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha\right)^{2}\right)^{2^{i}}+\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha+\left(\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha\right)^{2}
$$

Therefore,

$$
\left(\frac{\varphi_{2}}{\varphi_{3}}\right)^{2}=\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha+\left(\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha\right)^{2}
$$

or

$$
\begin{equation*}
\left(\frac{\varphi_{2}}{\varphi_{3}}\right)^{2}=\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha+\left(\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha\right)^{2}+1 \tag{19}
\end{equation*}
$$

If Eq. (19) holds, then

$$
\begin{equation*}
\left(1+\alpha^{2}\right)\left(\frac{\varphi_{2}}{\varphi_{3}}\right)^{2}+\alpha\left(\frac{\varphi_{2}}{\varphi_{3}}\right)+\alpha^{2}+\alpha+1=0 \tag{20}
\end{equation*}
$$

If $\alpha=1$, it is easy to obtain that $\beta=1$ from the definition of $\Gamma$. Moreover, $\varphi_{2}=0$ and thus $\operatorname{Tr}_{q}\left(\frac{\varphi_{2}}{\varphi_{3}}\right)=0$. In the following, we assume that $\alpha \neq 1$. Then after multiplying Eq. (20) by $\frac{\alpha^{2}+1}{\alpha^{2}}$ and simplifying, we get

$$
\left(\frac{\alpha^{2}+1}{\alpha} \cdot \frac{\varphi_{2}}{\varphi_{3}}\right)^{2}+\frac{\alpha^{2}+1}{\alpha} \cdot \frac{\varphi_{2}}{\varphi_{3}}=\left(\frac{\alpha^{2}+1}{\alpha}\right)^{2}+\frac{\alpha^{2}+1}{\alpha}
$$

and thus

$$
\frac{\varphi_{2}}{\varphi_{3}}=1 \quad \text { or } \frac{\alpha^{2}+\alpha+1}{\alpha^{2}+1} .
$$

It is clear that $\varphi_{2}+\varphi_{3} \neq 0 . \operatorname{By}\left(\frac{\varphi_{3}}{\varphi_{2}+\varphi_{3}}\right)^{2^{i}}=\frac{\varphi_{3}}{\varphi_{1}+\varphi_{3}}$, one has

$$
\left(\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}}\right)^{2^{i}-1}=\frac{\varphi_{1}+\varphi_{3}}{\varphi_{3}} \cdot \frac{\varphi_{3}}{\varphi_{2}+\varphi_{3}}=\frac{\varphi_{1}+\varphi_{3}}{\varphi_{2}+\varphi_{3}}=\frac{\alpha^{2^{i}+2}+\alpha^{2^{i}}+\alpha \beta}{\alpha^{2^{i+1}+1}+\alpha^{2^{i}} \beta+\alpha}=\frac{1}{\alpha^{2^{i}-1}} .
$$

Since $\operatorname{gcd}(i, n)=1$, one has $\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}}=\frac{1}{\alpha}$ and $\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}}=\frac{1}{\alpha}+1 \neq \frac{\alpha^{2}+\alpha+1}{\alpha^{2}+1}$.
Therefore, Eq. (19) does not hold and thus

$$
\begin{equation*}
\left(\frac{\varphi_{2}}{\varphi_{3}}\right)^{2}=\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha+\left(\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha\right)^{2} \tag{21}
\end{equation*}
$$

which implies $\operatorname{Tr}_{q}\left(\frac{\varphi_{2}}{\varphi_{3}}\right)=0$.

## 3 The permutation property of $V_{i}(x, y)$

In this section, we firstly give a general necessary and sufficient condition on the permutation property of the function $V_{i}$ from the closed butterfly. Throughout what follows, we always assume $n$ is an odd integer.

Recall that the univariate representation of $V_{i}$ have the following form

$$
\begin{equation*}
f(x)=\epsilon_{1} x^{q \cdot\left(2^{i}+1\right)}+\epsilon_{2} x^{q \cdot 2^{i}+1}+\epsilon_{3} x^{2^{i}+q}+\epsilon_{4} x^{2^{i}+1}, \quad \epsilon_{j} \in \mathbb{F}_{q} . \tag{22}
\end{equation*}
$$

Below we first present a necessary and sufficient conditions for $f(x)$ to be a permutation of $\mathbb{F}_{q^{2}}$ without imposing any additional restrictions on $\epsilon_{j}$. The following proposition investigates the permutation property of $f(x)$ defined by (22) over $\mathbb{F}_{q^{2}}$.

Proposition 15 Let $q=2^{n}$, $f(x)$ be defined by (22), $h(x)=\epsilon_{1} x^{2^{i}+1}+\epsilon_{2} x^{2^{i}}+\epsilon_{3} x+\epsilon_{4}$ and $g(x)=$ $x^{2^{i}+1} h(x)^{q-1}$. Define $\mu_{q+1}=\left\{x \in \mathbb{F}_{q^{2}}: x^{q+1}=1\right\}$ and

$$
T=\left\{\left(\frac{x y+1}{x+y}, \frac{x y}{x^{2}+y^{2}}\right) \quad: \quad x, y \in \mu_{q+1} \backslash\{1\}, y \neq x, x^{q}\right\} \subset \mathbb{F}_{q}^{2}
$$

Then $f(x)$ permutes $\mathbb{F}_{q^{2}}$ if and only if
(1) $\operatorname{gcd}\left(2^{i}+1, q-1\right)=1$;
(2) $h(x)=0$ has no solution in $\mu_{q+1}$;
(3) $g(x)=1$ if and only if $x=1$;
(4) there does not exist some $(X, Y) \in T$ such that the following equation holds:

$$
\begin{equation*}
\varphi_{1} X^{2^{i}}+\varphi_{2} X+\varphi_{3}\left(\sum_{j=0}^{i-1} Y^{2^{j}}\right)+\varphi_{4}=0 \tag{23}
\end{equation*}
$$

where $\varphi_{j}$ for $j=1,2,3,4$ are defined by (10).

Proof It is clear that $f(x)=x^{2^{i}+1} h\left(x^{q-1}\right)$. According to Lemma $9, f(x)$ permutes $\mathbb{F}_{q^{2}}$ if and only if $\operatorname{gcd}\left(2^{i}+1, q-1\right)=1$ and

$$
g(x)=x^{2^{i}+1} h(x)^{q-1}=\frac{\epsilon_{4} x^{2^{i}+1}+\epsilon_{3} x^{2^{i}}+\epsilon_{2} x+\epsilon_{1}}{\epsilon_{1} x^{2^{i}+1}+\epsilon_{2} x^{2^{i}}+\epsilon_{3} x+\epsilon_{4}}
$$

permutes $\mu_{q+1}$, which obviously implies that $h(x)=0$ has no solution in $\mu_{q+1}$ and $g(x)=1$ if and only if $x=1$. In the following, we assume that the conditions (1),(2) and (3) hold. Therefore, $g(x)$ permutes $\mu_{q+1}$ if and only if $g(x)+g(y)=0$ has no solution for $x, y \in \mu_{q+1} \backslash\{1\}$ with $x \neq y$. In fact, if $g(x)+g(y)=0$ for some $y=x^{q}$, then we have $g(x)=g(y)=g\left(x^{q}\right)=g(x)^{q}=g(x)^{-1}$ and thus $g(x)=1$, which means that $x=1$. Thus we can only consider the conditions such that $g(x)+g(y)=0$ has no solution for $x, y \in \mu_{q+1} \backslash\{1\}$ with $y \neq x, x^{q}$. Next, we prove the necessity and sufficiency of the condition (4).

The sufficiency of (4). Suppose $g(x)+g(y)=0$, i.e.,

$$
\frac{\epsilon_{4} x^{2^{i}+1}+\epsilon_{3} x^{2^{i}}+\epsilon_{2} x+\epsilon_{1}}{\epsilon_{1} x^{2^{i}+1}+\epsilon_{2} x^{2^{i}}+\epsilon_{3} x+\epsilon_{4}}=\frac{\epsilon_{4} y^{2^{i}+1}+\epsilon_{3} y^{2^{i}}+\epsilon_{2} y+\epsilon_{1}}{\epsilon_{1} y^{2^{i}+1}+\epsilon_{2} y^{2^{i}}+\epsilon_{3} y+\epsilon_{4}}
$$

After a routine calculation, we obtain

$$
\varphi_{1}(x+y)(x y+1)^{2^{i}}+\varphi_{2}(x+y)^{2^{i}}(x y+1)+\varphi_{4}(x+y)^{2^{i}+1}+\varphi_{3}\left(x^{2^{i}} y+x y^{2^{i}}\right)=0
$$

where $\varphi_{j}$ 's for $j=1,2,3,4$ are defined as in (10). By the previous discussion, we now only need to consider the case that $(x+y)(x y+1) \neq 0$. Therefore, the above equation is equivalent to

$$
\begin{equation*}
\varphi_{1}\left(\frac{x y+1}{x+y}\right)^{2^{i}}+\varphi_{2}\left(\frac{x y+1}{x+y}\right)+\varphi_{4}+\varphi_{3}\left(\frac{x^{2^{i}} y+x y^{2^{i}}}{(x+y)^{2^{i}+1}}\right)=0 \tag{24}
\end{equation*}
$$

Note that

$$
\frac{x^{2^{i}} y+x y^{2^{i}}}{(x+y)^{2^{2}+1}}=\frac{x^{2^{i}+1}+y y^{2^{i}+1}}{(x+y)^{2^{2}+1}}+1=\left(\frac{x}{x+y}\right)^{2^{i}+1}+\left(\frac{y}{x+y}\right)^{2^{i}+1}+1
$$

It follows from Lemma 12 (1) that the coefficient of $\varphi_{3}$ can be expressed in terms of Dickson polynomial as

$$
\frac{x^{2^{i}} y+x y^{2^{i}}}{(x+y)^{2^{i}+1}}=D_{2^{i}+1}\left(1, \frac{x y}{(x+y)^{2}}\right)+1
$$

In addition, by Lemma 12 (2), (4) and Lemma 13,

$$
\begin{equation*}
D_{2^{i}+1}(x, a)=D_{2^{i}}(x, a)+a D_{2^{i}-1}(x, a)=x^{2^{i}}+\sum_{j=0}^{i-1} a^{2^{j}} x^{2^{i}-2^{j+1}+1} \tag{25}
\end{equation*}
$$

Denote $X=\frac{x y+1}{x+y}$ and $Y=\frac{x y}{(x+y)^{2}}$. Then the coefficient of $\varphi_{3}$ can be written as

$$
D_{2^{i}+1}(1, Y)+1=1+\sum_{j=0}^{i-1} Y^{2^{j}}+1=\sum_{j=0}^{i-1} Y^{2^{j}}
$$

It is straightforward that $g(x)=g(y)$ can be rewritten as

$$
\begin{equation*}
\varphi_{1} X^{2^{i}}+\varphi_{2} X+\varphi_{3}\left(\sum_{j=0}^{i-1} Y^{2^{j}}\right)+\varphi_{4}=0 \tag{26}
\end{equation*}
$$

Thus, if there exist some $x, y \in \mu_{q+1}$ with $y \neq x, x^{q}$ such that $g(x)+g(y)=0$ holds, there must exist some $(X, Y) \in T$ such that Eq. (26) holds. Thus if the condition (4) holds, $g(x)$ permutes $\mu_{q+1}$.

The necessity of (4). On the contrary, if the condition (4) does not hold, which means that there exist some $(X, Y) \in T$ such that Eq. (26) holds, then there must exist some $x, y \in \mu_{q+1} \backslash\{1\}$ with $y \neq x, x^{q}$ such that $g(x)+g(y)=0$, which implies that $g(x)$ does not permute $\mu_{q+1}$.

On combining the sufficiency and necessity, we have proved the desired conclusion.

## Proof of the permutation part in Theorem 1.

In the following, we will prove the permutation part in Theorem 1 by verifying the conditions in Proposition 15.

First of all, if $\alpha=1$, it is easy to obtain that $\beta=1$ from the definition of $\Gamma$. In this case, the function

$$
f_{i}(x)=\left\{\begin{array}{l}
x^{q \cdot\left(2^{i}+1\right)}, \text { when } i \text { is even } \\
x^{2^{i}+q}, \text { when } i \text { is odd }
\end{array}\right.
$$

clearly permutes $\mathbb{F}_{q^{2}}$. In the following, we assume $\alpha \neq 1$ and will show the four items of Proposition 15 .
(1) Since $n$ is odd and $\operatorname{gcd}(i, n)=1$, we have $\operatorname{gcd}\left(2^{i}+1,2^{n}-1\right)=1$ due to the fact $\operatorname{gcd}\left(2^{i}+1,2^{n}-1\right) \mid$ $\operatorname{gcd}\left(2^{2 i}-1,2^{n}-1\right)=2^{\operatorname{gcd}(2 i, n)}-1=1$.
(2) Next we show that $h(x)=0$ has no solution in $\mu_{q+1} \backslash\{1\}\left(h(1)=\sqrt{\varphi_{3}} \neq 0\right.$ according to the definition). Suppose that there exists some $x_{0} \in \mu_{q+1} \backslash\{1\}$ satisfying

$$
\begin{equation*}
\epsilon_{1} x_{0}^{2^{i}+1}+\epsilon_{2} x_{0}^{2^{i}}+\epsilon_{3} x_{0}+\epsilon_{4}=0 \tag{27}
\end{equation*}
$$

Raising Eq. (27) to the $q$-th power and re-arranging it according to $x_{0}^{q}=x_{0}^{-1}$, we obtain

$$
\begin{equation*}
\epsilon_{4} x_{0}^{2^{i}+1}+\epsilon_{3} x_{0}^{2^{i}}+\epsilon_{2} x_{0}+\epsilon_{1}=0 \tag{28}
\end{equation*}
$$

Summing $\epsilon_{4} \times(27)$ and $\epsilon_{1} \times(28)$ gives

$$
\begin{equation*}
\varphi_{1} x_{0}^{2^{i}}+\varphi_{2} x_{0}+\varphi_{4}=0 \tag{29}
\end{equation*}
$$

Computing $\varphi_{4} \times(29)+\varphi_{1} \times(29)^{q} \times x_{0}^{2^{i}}$ yields

$$
\begin{equation*}
\varphi_{1} \varphi_{2} x_{0}^{x^{i}-1}+\varphi_{2} \varphi_{4} x_{0}+\varphi_{1}^{2}+\varphi_{4}^{2}=0 \tag{30}
\end{equation*}
$$

Furthermore, by computing (30) $\times x_{0}+(29) \times \varphi_{2}$, we obtain

$$
\begin{equation*}
\varphi_{2} \varphi_{4} x_{0}^{2}+\left(\varphi_{1}^{2}+\varphi_{2}^{2}+\varphi_{4}^{2}\right) x_{0}+\varphi_{2} \varphi_{4}=0 \tag{31}
\end{equation*}
$$

Note that in the above equation $\varphi_{2} \varphi_{4} \neq 0$. Otherwise, we have $\varphi_{1}^{2}+\varphi_{2}^{2}=\varphi_{4}^{2}$. Recall that $\varphi_{1}^{2}+\varphi_{2}^{2}=$ $\varphi_{4}\left(\varphi_{3}+\varphi_{4}\right)$ from (13) and (14). Thus we obtain $\varphi_{3} \varphi_{4}=0$, which is in contradiction with $\varphi_{3} \neq 0$ in the definition of $\Gamma$ and $\varphi_{4} \neq 0$ in Lemma 14 (1). Thus Eq. (31) becomes

$$
\begin{equation*}
x_{0}^{2}+\frac{\varphi_{1}^{2}+\varphi_{2}^{2}+\varphi_{4}^{2}}{\varphi_{2} \varphi_{4}} x_{0}+1=0 \tag{32}
\end{equation*}
$$

Note that

$$
\operatorname{Tr}_{q}\left(\frac{\varphi_{2} \varphi_{4}}{\varphi_{1}^{2}+\varphi_{2}^{2}+\varphi_{4}^{2}}\right)=\operatorname{Tr}_{q}\left(\frac{\varphi_{2} \varphi_{4}}{\varphi_{4} \varphi_{3}}\right)=\operatorname{Tr}_{q}\left(\frac{\varphi_{2}}{\varphi_{3}}\right)=0
$$

This implies that Eq. (32) has a solution $x_{0} \in \mathbb{F}_{q}$, which contradicts $\mu_{q+1} \backslash\{1\}$. Therefore, $h(x)=0$ has no solution in $\mu_{q+1}$.
(3) If there exists some $x_{0} \in \mu_{q+1} \backslash\{1\}$ such that $g\left(x_{0}\right)=1$, then we have

$$
\begin{equation*}
\left(\epsilon_{1}+\epsilon_{4}\right) x_{0}^{2^{i}+1}+\left(\epsilon_{2}+\epsilon_{3}\right) x_{0}^{2^{i}}+\left(\epsilon_{2}+\epsilon_{3}\right) x_{0}+\epsilon_{1}+\epsilon_{4}=0 \tag{33}
\end{equation*}
$$

According to Lemma 10 , we know that for any $x_{0} \in \mu_{q+1} \backslash\{1\}$, there exists a unique element $y_{0} \in$ $\mathbb{F}_{q}$ such that $x_{0}=\frac{y_{0}+\gamma}{y_{0}+\gamma^{2}}$, where $\gamma \in \mathbb{F}_{2^{2}} \backslash \mathbb{F}_{2}$. By plugging $x_{0}=\frac{y_{0}+\gamma}{y_{0}+\gamma^{2}}$ into Eq. (33) and a routine rearrangement, we obtain

$$
\begin{equation*}
y_{0}^{2^{i}}+y_{0}+\frac{\varepsilon_{1}+\varepsilon_{4}}{\epsilon_{1}+\epsilon_{2}+\epsilon_{3}+\epsilon_{4}}=0 \tag{34}
\end{equation*}
$$

where $\varepsilon_{1}, \varepsilon_{4}$ are defined as in (9) satisfying that $\varepsilon_{1}+\varepsilon_{4}=\epsilon_{1}+\epsilon_{4}$ for even $i$ and $\varepsilon_{1}+\varepsilon_{4}=\epsilon_{2}+\epsilon_{3}$ for odd $i$. In other words, $\varepsilon_{1}+\varepsilon_{4}$ corresponds to $\sqrt{\varphi_{4}}$ for even $i$ and $\sqrt{\varphi_{3}+\varphi_{4}}$ for odd $i$. By Lemma 14 (2) and (3), we have

$$
\operatorname{Tr}_{q}\left(\frac{\varepsilon_{1}+\varepsilon_{4}}{\epsilon_{1}+\epsilon_{2}+\epsilon_{3}+\epsilon_{4}}\right)=1
$$

This implies (34) has no solution in $\mathbb{F}_{q}$. Hence $g(x)=1$ if and only if $x=1$.
(4) Recall that $Y=\frac{x y}{x^{2}+y^{2}}$ for some $x, y \in \mu_{q+1} \backslash\{1\}$ with $x \neq y$. Note that $\frac{y}{x+y} \in \mathbb{F}_{q^{2}} \backslash \mathbb{F}_{q}$ is a solution to the equation $z^{2}+z+Y=0$. This implies $\operatorname{Tr}_{q}(Y)=1$. It is clear that Eq. (26) required in Proposition 15 is equivalent to

$$
\sum_{j=0}^{i-1} Y^{2^{j}}=\frac{\varphi_{1}}{\varphi_{3}} X^{2^{i}}+\frac{\varphi_{2}}{\varphi_{3}} X+\frac{\varphi_{4}}{\varphi_{3}}=\left(\frac{\varphi_{2}}{\varphi_{3}} X\right)^{2^{i}}+\frac{\varphi_{2}}{\varphi_{3}} X+\frac{\varphi_{4}}{\varphi_{3}}
$$

By $\operatorname{Tr}_{q}(Y)=1$ we have

$$
\operatorname{Tr}_{q}\left(\sum_{j=0}^{i-1} Y^{2^{j}}\right)=\left\{\begin{array}{l}
0, \text { when } i \text { is even }  \tag{35}\\
1, \text { when } i \text { is odd }
\end{array}\right.
$$

on the other hand, the expression on the right hand side satisfies

$$
\operatorname{Tr}_{q}\left(\left(\frac{\varphi_{2}}{\varphi_{3}} X\right)^{2^{i}}+\frac{\varphi_{2}}{\varphi_{3}} X+\frac{\varphi_{4}}{\varphi_{3}}\right)=\left\{\begin{array}{l}
1, \text { when } i \text { is even } \\
0, \text { when } i \text { is odd }
\end{array}\right.
$$

according to Lemma 14. It is clear that Eq. (26) does not hold for any $X, Y \in \mathbb{F}_{q}$.
Up to now, all the four items in Proposition 15 are confirmed. Hence the function $V_{i}(x, y)$ in Theorem 1 permutes $\mathbb{F}_{q}^{2}$.

## 4 The boomerang uniformity of $V_{i}(x, y)$

In this section, we will prove that the function

$$
V_{i}:=\left(R_{i}(x, y), R_{i}(y, x)\right)
$$

with $R_{i}(x, y)=(x+\alpha y)^{2^{i}+1}+\beta y^{2^{i}+1}$ has boomerang uniformity 4 when the pair $(\alpha, \beta)$ is taken from the set $\Gamma$ as in given in Theorem 1. Here and hereafter, we assume that $n$ is odd, $q=2^{n}$ and $(\alpha, \beta) \in \Gamma$.

First of all, the condition $\beta \neq(\alpha+1)^{2^{i}+1}$ in Lemma 8 corresponds to the condition $\varphi_{3} \neq 0$ in $\Gamma$. Hence the differential uniformity of $V_{i}$ with $R_{i}(x, y)=(x+\alpha y)^{2^{i}+1}+\beta y^{2^{i}+1}$ is at most 4 for any $(\alpha, \beta) \in \Gamma$. Furthermore, Canteaut, Perrin and Tian [5] showed that if $V_{i}$ is APN then it operates on 6 bits. Therefore, the differential uniformity of $V_{i}$ is equal to 4 in other cases. Since $V_{i}$ in Theorem 1 permutes $\mathbb{F}_{q}^{2}$ and has differential uniformity 4 , we can use Lemma 5 to show the boomerang uniformity of $V_{i}$. For any $(a, b) \in \mathbb{F}_{q}^{2}$, denote

$$
S_{V_{i},(a, b)}(x, y)=V_{i}(x+a, y+b)+V_{i}(x, y)+V_{i}(a, b)
$$

and

$$
\operatorname{Im}_{V_{i},(a, b)}=\left\{S_{V_{i},(a, b)}(x, y):(x, y) \in \mathbb{F}_{q}^{2}\right\}
$$

According to Lemma 5 , we need to determine $\left(a_{1}, b_{1}\right),\left(a_{2}, b_{2}\right) \in \mathbb{F}_{q}^{2} \backslash\{(0,0)\}$ satisfying $S_{V_{i},\left(a_{1}, b_{1}\right)}\left(a_{2}, b_{2}\right)=$ $(0,0)$, and then to prove that for any such pairs the equation $\operatorname{Im}_{V_{i},\left(a_{1}, b_{1}\right)}=\operatorname{Im}_{V_{i},\left(a_{2}, b_{2}\right)}$ holds.
4.1 The solutions of $S_{V_{i},\left(a_{1}, b_{1}\right)}\left(a_{2}, b_{2}\right)=(0,0)$

The solution of the equation $S_{V_{i},\left(a_{1}, b_{1}\right)}\left(a_{2}, b_{2}\right)=(0,0)$ is studied in the following proposition.
Proposition 16 Let $V_{i}$ be defined as in Theorem 1 with $(\alpha, \beta) \in \Gamma$ and $\varphi_{j}$ 's for $j=1,2,3,4$ be defined as in (10). Then the elements $\left(a_{1}, b_{1}\right),\left(a_{2}, b_{2}\right) \in \mathbb{F}_{q}^{2} \backslash\{(0,0)\}$ such that

$$
V_{i}\left(a_{1}+a_{2}, b_{1}+b_{2}\right)+V_{i}\left(a_{1}, b_{1}\right)+V_{i}\left(a_{2}, b_{2}\right)=(0,0)
$$

satisfy $\left(a_{2}, b_{2}\right)=X \cdot\left(a_{1}, b_{1}\right)$, where $X$ is a $2 \times 2$ matrix taken from the following set

$$
\left\{\left[\begin{array}{ll}
1 & 0  \tag{36}\\
0 & 1
\end{array}\right],\left[\begin{array}{cc}
\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha+1 & \frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \\
\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} & \frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha
\end{array}\right],\left[\begin{array}{cc}
\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha & \frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \\
\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} & \frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha+1
\end{array}\right]\right\}
$$

Proof Note that the equation

$$
S_{V_{i},\left(a_{1}, b_{1}\right)}\left(a_{2}, b_{2}\right)=V_{i}\left(a_{1}+a_{2}, b_{1}+b_{2}\right)+V_{i}\left(a_{1}, b_{1}\right)+V_{i}\left(a_{2}, b_{2}\right)=(0,0)
$$

can be rewritten as

$$
\left\{\begin{array}{l}
\left(a_{1}+\alpha b_{1}\right) a_{2}^{2^{i}}+\left(a_{1}^{2^{i}}+\alpha^{2^{i}} b_{1}^{2^{i}}\right) a_{2}+\left({\alpha^{2^{i}}}^{i} a_{1}+\left(\alpha^{2^{i}+1}+\beta\right) b_{1}\right) b_{2}^{2^{i}}+\left(\alpha a_{1}^{2^{i}}+\left(\alpha^{2^{i}+1}+\beta\right) b_{1}^{2^{i}}\right) b_{2}=0  \tag{37.1}\\
\left(\left(\alpha^{2^{i}+1}+\beta\right) a_{1}+{\left.\alpha^{2^{i}} b_{1}\right) a_{2}^{2^{i}}+\left(\left(\alpha^{2^{i}+1}+\beta\right) a_{1}^{2^{i}}+\alpha b_{1}^{2^{i}}\right) a_{2}+\left(\alpha a_{1}+b_{1}\right) b_{2}^{i^{i}}+\left(\alpha^{2^{i}} a_{1}^{2^{i}}+b_{1}^{2^{i}}\right) b_{2}=0 .}^{\text {and }} .\right.
\end{array}\right.
$$

Let $\varphi_{j}$ 's for $j=1,2,3,4$ be defined as in (10). Eliminating the terms $a_{2}^{2^{i}}$ in the above equations by computing $(37.1) \times\left(\left(\alpha^{2^{i}+1}+\beta\right) a_{1}+\alpha^{2^{i}} b_{1}\right)+(37.2) \times\left(a_{1}+\alpha b_{1}\right)$, we obtain

$$
\begin{equation*}
\lambda_{1} a_{2}+\lambda_{2} b_{2}^{2^{i}}+\lambda_{3} b_{2}=0, \tag{38}
\end{equation*}
$$

where the coefficients are given by

$$
\left\{\begin{array}{l}
\lambda_{1}=\left(\varphi_{1}+\varphi_{3}\right) a_{1}^{2^{i}} b_{1}+\left(\varphi_{2}+\varphi_{3}\right) a_{1} b_{1}^{2^{i}}+\left(\varphi_{3}+\varphi_{4}\right) b_{1}^{2^{i}+1} \\
\lambda_{2}=\left(\varphi_{2}+\varphi_{3}\right) a_{1}^{2}+\varphi_{3} a_{1} b_{1}+\left(\varphi_{2}+\varphi_{3}\right) b_{1}^{2} \\
\lambda_{3}=\left(\varphi_{1}+\varphi_{3}\right) a_{1}^{2^{i}+1}+\varphi_{4} a_{1} b_{1}^{2^{i}}+\left(\varphi_{2}+\varphi_{3}\right) b_{1}^{2^{i}+1}
\end{array}\right.
$$

$\varphi_{1}, \varphi_{2}, \varphi_{3}$ are defined as in (10) and $\varphi_{4}$ corresponds to the case that $i$ is even. Hereafter, we assume $\varphi_{4}$ is restricted to the case of even $i$, i.e, $\varphi_{4}=\left(\alpha^{2^{i}+1}+\beta+1\right)^{2}$.

When $b_{1}=0$, we have $a_{1} \neq 0, \lambda_{1}=0, \lambda_{2}=\left(\varphi_{2}+\varphi_{3}\right) a_{1}^{2}$ and $\lambda_{3}=\left(\varphi_{1}+\varphi_{3}\right) a_{1}^{2^{i}+1}$. Moreover, Eq. (38) becomes $\lambda_{2} b_{2}^{2^{i}}=\lambda_{3} b_{2}$. This together with Lemma 14 (1) implies

$$
b_{2}=0 \text { or } b_{2}=\left(\frac{\varphi_{1}+\varphi_{3}}{\varphi_{2}+\varphi_{3}}\right)^{\frac{1}{2^{i}-1}} a_{1}=\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} a_{1} .
$$

Note that in the case of $b_{1}=0$, Eq. (37.1) becomes

$$
\left(\frac{a_{2}}{a_{1}}\right)^{2^{i}}+\frac{a_{2}}{a_{1}}=\left(\frac{\alpha b_{2}}{a_{1}}\right)^{2^{i}}+\frac{\alpha b_{2}}{a_{1}} .
$$

Therefore, if $b_{2}=0$, then $a_{2}=a_{1}$; if $b_{2}=\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} a_{1}$, then $a_{2}=\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha a_{1}$ or $a_{2}=\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha a_{1}+a_{1}$.
When $b_{1} \neq 0$, after eliminating the terms $b_{2}^{2^{i}}$ by computing $(37.1) \times\left(\left(\alpha^{2^{i}+1}+\beta\right) a_{1}^{2^{i}}+\alpha b_{1}^{2^{i}}\right)+(37.2) \times$ $\left(a_{1}^{2^{i}}+\alpha^{2^{i}} b_{1}^{2^{i}}\right)$, we obtain

$$
\begin{equation*}
\eta_{1} a_{2}^{2^{i}}+\eta_{2} b_{2}^{2^{i}}+\eta_{3} b_{2}=0 \tag{39}
\end{equation*}
$$

where

$$
\left\{\begin{array}{l}
\eta_{1}=\lambda_{1} \\
\eta_{2}=\left(\varphi_{2}+\varphi_{3}\right) a_{1}^{2^{i}+1}+\varphi_{4} a_{1}^{2^{i}} b_{1}+\left(\varphi_{1}+\varphi_{3}\right) b_{1}^{2^{i}+1} \\
\eta_{3}=\left(\varphi_{1}+\varphi_{3}\right) a_{1}^{2^{i+1}}+\varphi_{3} a_{1}^{2^{i}} b_{1}^{2^{i}}+\left(\varphi_{1}+\varphi_{3}\right) b_{1}^{2^{i+1}}
\end{array}\right.
$$

Furthermore, computing $(38)^{2^{i}}+\lambda_{1}^{2^{i}-1} \times(39)$, we eliminate the term ${a_{2}^{2^{i}}}$ and obtain

$$
\begin{equation*}
\lambda_{2}^{2^{i}} b_{2}^{2^{2 i}-1}+\left(\lambda_{1}^{2^{i}-1} \eta_{2}+\lambda_{3}^{2^{i}}\right) b_{2}^{2^{i}-1}+\lambda_{1}^{2^{i}-1} \eta_{3}=0 \tag{40}
\end{equation*}
$$

Here we note that $\lambda_{2} \neq 0$. Otherwise one has $\left(\varphi_{2}+\varphi_{3}\right) a_{1}^{2}+\varphi_{3} a_{1} b_{1}+\left(\varphi_{2}+\varphi_{3}\right) b_{1}^{2}=0$, i.e.,

$$
\left(\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \cdot \frac{a_{1}}{b_{1}}\right)^{2}+\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \cdot \frac{a_{1}}{b_{1}}+\left(\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}}\right)^{2}=0
$$

which is in contradiction with the fact $\operatorname{Tr}_{q}\left(\frac{\varphi_{2}}{\varphi_{3}}\right)=0$ by Lemma 14 (4).

In addition, since the differential uniformity of $V_{i}$ is 4, Eq. (40) has three nonzero solutions $b_{2}=b_{1}, \bar{b}$ and $\bar{b}+b_{1}$ and we only need to obtain the expression of $\bar{b}$. Clearly, $\tilde{b}_{2}=b_{1}^{2^{i}-1}$ is a solution of

$$
\begin{equation*}
\lambda_{2}^{2^{i}} \tilde{b}_{2}^{2^{i}+1}+\left(\lambda_{1}^{2^{i}-1} \eta_{2}+\lambda_{3}^{2^{i}}\right) \tilde{b}_{2}+\lambda_{1}^{2^{i}-1} \eta_{3}=0 \tag{41}
\end{equation*}
$$

Hence, Eq. (41) can be written as

$$
\lambda_{2}^{2^{i}}\left(\tilde{b}_{2}+b_{1}^{2^{i}-1}\right)\left(\tilde{b}_{2}^{2^{i}}+b_{1}^{2^{i}-1} \tilde{b}_{2}^{2^{i}-1}+b_{1}^{2 \cdot\left(2^{i}-1\right)} \tilde{b}_{2}^{2^{i}-2}+\cdots+b_{1}^{\left(2^{i}-1\right) \cdot\left(2^{i}-1\right)} \tilde{b}_{2}+c\right)=0
$$

where $c=\frac{\lambda_{1}^{2^{i}-1} \eta_{3}}{\lambda_{2}^{\lambda^{i} b_{1}^{2 i}-1}}$. Now we consider the equation

$$
\begin{equation*}
\tilde{b}_{2}^{2^{i}}+b_{1}^{2^{i}-1} \tilde{b}_{2}^{2^{i}-1}+b_{1}^{2 \cdot\left(2^{i}-1\right)} \tilde{b}_{2}^{2^{i}-2}+\cdots+b_{1}^{\left(2^{i}-1\right) \cdot\left(2^{i}-1\right)} \tilde{b}_{2}+c=0 \tag{42}
\end{equation*}
$$

Let $\hat{b}_{2}=\frac{1}{\hat{b}_{2}+b_{1}^{2^{i}-1}}$. Then Eq. (42) becomes

$$
\hat{b}_{2}^{2^{i}}+\frac{b_{1}^{2^{i}-1}}{c} \hat{b}_{2}+\frac{1}{c}=0
$$

i.e.,

$$
\begin{equation*}
\left(\frac{c^{\frac{1}{2^{i}-1}}}{b_{1}} \hat{b}_{2}\right)^{2^{i}}+\frac{c^{\frac{1}{2^{i}-1}}}{b_{1}} \hat{b}_{2}+\frac{c^{\frac{1}{2^{i}-1}}}{b_{1}^{2^{i}}}=0 \tag{43}
\end{equation*}
$$

In addition, we have

$$
\begin{aligned}
c^{\frac{1}{2^{i}-1}} & =\left(\frac{\lambda_{1}^{2^{i}-1} \eta_{3}}{\lambda_{2}^{2^{i}} b_{1}^{2^{i}-1}}\right)^{\frac{1}{2^{i}-1}} \\
& =\frac{\lambda_{1}}{b_{1}}\left(\frac{\left(\varphi_{1}+\varphi_{3}\right)\left(a_{1}^{2^{i+1}}+\frac{\varphi_{3}}{\varphi_{1}+\varphi_{3}} a_{1}^{2^{i}} b_{1}^{2^{i}}+b_{1}^{2^{i+1}}\right)}{\left(\varphi_{2}+\varphi_{3}\right)^{2^{i}}\left(a_{1}^{2^{i+1}}+\frac{\varphi_{3}^{2 i}}{\left(\varphi_{2}+\varphi_{3}\right)^{2^{2}}} a_{1}^{2^{i}} b_{1}^{2^{i}}+b_{1}^{2^{i+1}}\right)}\right)^{\frac{1}{2^{i}-1}} \\
& =\frac{\lambda_{1}}{b_{1}}\left(\frac{\varphi_{1}+\varphi_{3}}{\left(\varphi_{2}+\varphi_{3}\right)^{2^{i}}}\right)^{\frac{1}{2^{i}-1}} \\
& =\frac{\lambda_{1}}{b_{1} \varphi_{3}}
\end{aligned}
$$

where the last two equalities follow from Lemma 14 (1). Moreover,

$$
\begin{aligned}
\frac{c^{\frac{1}{2^{i}-1}}}{b_{1}^{2^{i}}} & =\frac{\lambda_{1}}{\varphi_{3} b_{1}^{2^{i}+1}} \\
& =\frac{\left(\varphi_{1}+\varphi_{3}\right) a_{1}^{2^{i}}+\left(\varphi_{2}+\varphi_{3}\right) a_{1} b_{1}^{2^{i}-1}+\left(\varphi_{3}+\varphi_{4}\right) b_{1}^{2^{i}}}{\varphi_{3} b_{1}^{2^{i}}} \\
& =\left(\frac{\left(\varphi_{2}+\varphi_{3}\right) a_{1}}{\varphi_{3} b_{1}}\right)^{2^{i}}+\frac{\left(\varphi_{2}+\varphi_{3}\right) a_{1}}{\varphi_{3} b_{1}}+\frac{\varphi_{3}+\varphi_{4}}{\varphi_{3}} \\
& =\left(\frac{\left(\varphi_{2}+\varphi_{3}\right) a_{1}}{\varphi_{3} b_{1}}+u\right)^{2^{i}}+\frac{\left(\varphi_{2}+\varphi_{3}\right) a_{1}}{\varphi_{3} b_{1}}+u,
\end{aligned}
$$

where $u=\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha$ due to Lemma 14 (2). Hence, from Eq. (43), we have

$$
\frac{c^{\frac{1}{2^{i}-1}}}{b_{1}} \hat{b}_{2} \in\left\{\frac{\left(\varphi_{2}+\varphi_{3}\right) a_{1}}{\varphi_{3} b_{1}}+u, \frac{\left(\varphi_{2}+\varphi_{3}\right) a_{1}}{\varphi_{3} b_{1}}+u+1\right\}
$$

which means that there are exactly two solutions in $\mathbb{F}_{q}$ for Eq. (42). W.L.O.G., we only consider the first expression here. Namely, we get

$$
\hat{b}_{2}=\frac{b_{1}}{c^{\frac{1}{2^{i}-1}}}\left(\frac{\left(\varphi_{2}+\varphi_{3}\right) a_{1}}{\varphi_{3} b_{1}}+u\right)=\frac{\left(\varphi_{2}+\varphi_{3}\right) a_{1} b_{1}+\varphi_{3} u b_{1}^{2}}{\lambda_{1}} .
$$

Thus,

$$
\tilde{b}_{2}=\frac{1}{\hat{b}_{2}}+b_{1}^{2^{i}-1}=\frac{\lambda_{1}}{\left(\varphi_{2}+\varphi_{3}\right) a_{1} b_{1}+\varphi_{3} u b_{1}^{2}}+b_{1}^{2^{i}-1}
$$

is one solution of Eq. (42). Furthermore, one solution of Eq. (40) is

$$
\begin{aligned}
b_{2} & =\left(\tilde{b}_{2}\right)^{\frac{1}{2^{i}-1}} \\
& =\left(\frac{\left(\varphi_{1}+\varphi_{3}\right) a_{1}^{2^{i}}+\varphi_{3} u^{2^{i}} b_{1}^{2^{i}}}{\left(\varphi_{2}+\varphi_{3}\right) a_{1}+\varphi_{3} u b_{1}}\right)^{\frac{1}{2^{i}-1}} \\
& =\left(\frac{\varphi_{1}+\varphi_{3}}{\varphi_{2}+\varphi_{3}}\right)^{\frac{1}{2^{i}-1}} \cdot\left(\frac{a_{1}^{2^{i}}+\frac{\varphi_{3}}{\varphi_{1}+\varphi_{3}} u^{2^{i}} b_{1}^{2^{i}}}{a_{1}+\frac{\varphi_{3}}{\varphi_{2}+\varphi_{3}} u b_{1}}\right)^{\frac{1}{2^{i}-1}} \\
& =\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}}\left(a_{1}+\frac{\varphi_{3}}{\varphi_{2}+\varphi_{3}} u b_{1}\right)(\text { by Lemma } 14(1)) \\
& =\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} a_{1}+\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha b_{1}\left(\text { recall that } u=\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha\right) .
\end{aligned}
$$

It follows directly from Eq. (38) that

$$
a_{2}=\frac{\lambda_{2}}{\lambda_{1}} b_{2}^{2^{i}}+\frac{\lambda_{3}}{\lambda_{1}} b_{2}=\left(\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha+1\right) a_{1}+\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} b_{1} .
$$

4.2 The proof of $\operatorname{Im}_{V_{i},\left(a_{1}, b_{1}\right)}=\operatorname{Im}_{V_{i},\left(a_{2}, b_{2}\right)}$

In this subsection, we prove that for any $\left(a_{1}, b_{1}\right),\left(a_{2}, b_{2}\right) \in \mathbb{F}_{q}^{2} \backslash\{(0,0)\}$ satisfying $S_{V_{i},\left(a_{1}, b_{1}\right)}\left(a_{2}, b_{2}\right)=$ $(0,0), \operatorname{Im}_{V_{i},\left(a_{1}, b_{1}\right)}=\operatorname{Im}_{V_{i},\left(a_{2}, b_{2}\right)}$.

According to Eq. (37), we know that for any $\left(a_{1}, b_{1}\right) \in \mathbb{F}_{q}^{2}, S_{V_{i},\left(a_{1}, b_{1}\right)}(x, y)$ can be represented as

$$
S_{V_{i},\left(a_{1}, b_{1}\right)}(x, y)=A_{1}\left[\begin{array}{c}
x^{2^{i}} \\
x
\end{array}\right]+B_{1}\left[\begin{array}{c}
y^{2^{i}} \\
y
\end{array}\right]
$$

where

$$
A_{1}=\left[\begin{array}{cc}
a_{1}+\alpha b_{1}, & a_{1}^{2^{i}}+\alpha^{2^{i}} b_{1}^{2^{i}} \\
\left(\alpha^{2^{i}+1}+\beta\right) a_{1}+\alpha^{2^{i}} b_{1},\left(\alpha^{2^{i}+1}+\beta\right) a_{1}^{2^{i}}+\alpha b_{1}^{2^{i}}
\end{array}\right] \triangleq\left[\begin{array}{l}
a_{11}, a_{12} \\
a_{13}, a_{14}
\end{array}\right]
$$

and

$$
B_{1}=\left[\begin{array}{cc}
\alpha^{2^{i}} a_{1}+\left(\alpha^{2^{i}+1}+\beta\right) b_{1}, \alpha a_{1}^{2^{i}}+\left(\alpha^{2^{i}+1}+\beta\right) b_{1}^{2^{i}} \\
\alpha a_{1}+b_{1}, & \alpha^{2^{i}} a_{1}^{2^{i}}+b_{1}^{2^{i}}
\end{array}\right] \triangleq\left[\begin{array}{l}
b_{11}, b_{12} \\
b_{13}, b_{14}
\end{array}\right] .
$$

For the three relations between $\left(a_{1}, b_{1}\right),\left(a_{2}, b_{2}\right) \in \mathbb{F}_{q}^{2} \backslash\{(0,0)\}$ presented in Proposition 16 such that $S_{V_{i},\left(a_{1}, b_{1}\right)\left(a_{2}, b_{2}\right)}=(0,0)$, it is clear that if $a_{2}=a_{1}$ and $b_{2}=b_{1}$, we have $\operatorname{Im}_{V_{i},\left(a_{1}, b_{1}\right)}=\operatorname{Im}_{V_{i},\left(a_{2}, b_{2}\right)}$. In addition, if we have proved that $\operatorname{Im}_{V_{i},\left(a_{1}, b_{1}\right)}=\operatorname{Im}_{V_{i},\left(a_{2}, b_{2}\right)}$ holds for the second relation in Proposition 16 , then so does it for the third relation since the sum of two same subspace equals to the subspace. Therefore, it suffices to show that $\operatorname{Im}_{V_{i},\left(a_{1}, b_{1}\right)}=\operatorname{Im}_{V_{i},\left(a_{2}, b_{2}\right)}$ holds for the second relation in Proposition 16. Below we will again restrict $\varphi_{4}$ to the case of even $i$.

Let $u=\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha$. Then $u^{2^{i}}=u+\frac{\varphi_{3}+\varphi_{4}}{\varphi_{3}}$. Moreover, $a_{2}=(u+1) a_{1}+\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} b_{1}$ and $b_{2}=\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} a_{1}+u b_{1}$. Furthermore, we get

$$
a_{2}^{2^{i}}=\left(u^{2^{i}}+1\right) a_{1}^{2^{i}}+\left(\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}}\right)^{2^{i}} b_{1}^{2^{i}}=\left(u+\frac{\varphi_{4}}{\varphi_{3}}\right) a_{1}^{2^{i}}+\frac{\varphi_{1}+\varphi_{3}}{\varphi_{3}} b_{1}^{2^{i}}
$$

and

$$
b_{2}^{2^{i}}=\left(\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}}\right)^{2^{i}} a_{1}^{2^{i}}+u^{2^{i}} b_{1}^{2^{i}}=\frac{\varphi_{1}+\varphi_{3}}{\varphi_{3}} a_{1}^{2^{i}}+\left(u+\frac{\varphi_{3}+\varphi_{4}}{\varphi_{3}}\right) b_{1}^{2^{i}}
$$

Therefore, in $S_{V_{i},\left(a_{2}, b_{2}\right)}(x, y)$, we have

$$
A_{2}=\left[\begin{array}{cc}
a_{2}+\alpha b_{2}, & a_{2}^{2^{i}}+\alpha^{2^{i}} b_{2}^{2^{i}} \\
\left(\alpha^{2^{i}+1}+\beta\right) a_{2}+\alpha^{2^{i}} b_{2},\left(\alpha^{2^{i}+1}+\beta\right) a_{2}^{2^{i}}+\alpha b_{2}^{2^{i}}
\end{array}\right] \triangleq\left[\begin{array}{l}
a_{21}, a_{22} \\
a_{23}, a_{24}
\end{array}\right]
$$

and

$$
B_{2}=\left[\begin{array}{cc}
\alpha^{2^{i}} a_{2}+\left(\alpha^{2^{i}+1}+\beta\right) b_{2}, \alpha a_{2}^{2^{i}}+\left(\alpha^{2^{i}+1}+\beta\right) b_{2}^{2^{i}} \\
\alpha a_{2}+b_{2}, & \alpha^{2^{i}} a_{2}^{2^{i}}+b_{2}^{2^{i}}
\end{array}\right] \triangleq\left[\begin{array}{l}
b_{21}, b_{22} \\
b_{23}, b_{24}
\end{array}\right],
$$

where the explicit expressions of entries in $A_{2}$ and $B_{2}$ in terms of $a_{1}, b_{1}$ are given as follows:

$$
\begin{aligned}
a_{21}= & a_{2}+\alpha b_{2} \\
= & (u+1) a_{1}+\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} b_{1}+\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha a_{1}+u \alpha b_{1} \\
= & \left(u+1+\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha\right) a_{1}+\left(\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}}+u \alpha\right) b_{1} \\
= & \left.a_{1}+\left(\alpha^{2}+1\right) \frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} b_{1} \text { (recall that } u=\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha\right), \\
a_{22}= & a_{2}^{2^{i}}+\alpha^{2^{i}} b_{2}^{2^{i}} \\
= & a_{1}^{2^{i}}+\left(\alpha^{2^{i+1}}+1\right) \frac{\varphi_{1}+\varphi_{3}}{\varphi_{3}} b_{1}^{2^{i}}(\text { due to Lemma } 14(1)), \\
a_{23}= & \left(\alpha^{2^{i}+1}+\beta\right) a_{2}+\alpha^{2^{i}} b_{2} \\
= & \left(\left(\alpha^{2^{i}+1}+\beta\right)(u+1)+\alpha^{2^{i}} \frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}}\right) a_{1}+\left(\left(\alpha^{2^{i}+1}+\beta\right) \frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}}+\alpha^{2^{i}} u\right) b_{1} \\
= & \left(\frac{\left(\varphi_{2}+\varphi_{3}\right)\left(\varphi_{1}+\varphi_{3}\right)}{\varphi_{3}}+\alpha^{2^{i}+1}+\beta\right) a_{1}+\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \beta b_{1}, \\
a_{24}= & \left(\alpha^{2^{i}+1}+\beta\right) a_{2}^{2^{i}}+\alpha b_{2}^{2^{i}} \\
= & \left(\left(\alpha^{2^{i}+1}+\beta\right)\left(u+\frac{\varphi_{4}}{\varphi_{3}}\right)+\alpha \frac{\varphi_{1}+\varphi_{3}}{\varphi_{3}}\right) a_{1}^{2^{i}} \\
& +\left(\left(\alpha^{2^{i}+1}+\beta\right) \frac{\varphi_{1}+\varphi_{3}}{\varphi_{3}}+\alpha\left(u+\frac{\varphi_{3}+\varphi_{4}}{\varphi_{3}}\right)\right) b_{1}^{2^{i}} \\
= & \left(\frac{\left(\varphi_{2}+\varphi_{3}\right)\left(\varphi_{1}+\varphi_{3}\right)}{\varphi_{3}}+\alpha^{2^{i}+1}+\beta\right) a_{1}^{2^{i}}+\frac{\varphi_{1}+\varphi_{3}}{\varphi_{3}} \beta b_{1}^{2^{i}}(\text { due to }(15) \text { and }(16)), \\
& \\
b_{21}= & \alpha^{2^{i}} a_{2}+\left(\alpha^{2^{i}+1}+\beta\right) b_{2}=\left(\alpha^{2^{i}}+\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \beta\right) a_{1}+\frac{\left(\varphi_{2}+\varphi_{3}\right)\left(\varphi_{1}+\varphi_{3}\right)}{\varphi_{3}} b_{1}, \\
b_{22}= & \alpha a_{2}^{2^{i}+\left(\alpha^{2^{i}+1}+\beta\right) b_{2}^{2^{i}}=\left(\alpha+\frac{\varphi_{1}+\varphi_{3}}{\varphi_{3}} \beta\right) a_{1}^{2^{i}}+\frac{\left(\varphi_{2}+\varphi_{3}\right)\left(\varphi_{1}+\varphi_{3}\right)}{\varphi_{3}} b_{1}^{2^{i}},} \\
b_{23}= & \alpha a_{2}+b_{2}=\left(\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}}\left(\alpha^{2}+1\right)+\alpha\right) a_{1}, \\
b_{24}= & \alpha^{2^{i}} a_{2}^{2^{i}}+b_{2}^{2^{i}}=\left(\frac{\varphi_{1}+\varphi_{3}}{\varphi_{3}}\left(\alpha^{2^{i+1}}+1\right)+\alpha^{2^{i}}\right) a_{1}^{2^{i}} .
\end{aligned}
$$

Note that the determinants of $A_{1}$ and $B_{1}$ are

$$
\begin{aligned}
\operatorname{Det}\left(A_{1}\right) & =a_{11} a_{14}+a_{12} a_{13} \\
& =\left(\varphi_{1}+\varphi_{3}\right) a_{1}^{2^{i}} b_{1}+\left(\varphi_{2}+\varphi_{3}\right) a_{1} b_{1}^{2^{i}}+\left(\varphi_{3}+\varphi_{4}\right) b_{1}^{2^{i}+1}
\end{aligned}
$$

and

$$
\begin{aligned}
\operatorname{Det}\left(B_{1}\right) & =b_{11} b_{14}+b_{12} b_{13} \\
& =\left(\varphi_{3}+\varphi_{4}\right) a_{1}^{2^{i}+1}+\left(\varphi_{2}+\varphi_{3}\right) a_{1}^{2^{i}} b_{1}+\left(\varphi_{1}+\varphi_{3}\right) a_{1} b_{1}^{2^{i}}
\end{aligned}
$$

Now we consider the necessary and sufficient conditions such that $\operatorname{Det}\left(A_{1}\right)=0$. Clearly, from $\operatorname{Det}\left(A_{1}\right)=0$, we have $b_{1}=0$ or

$$
\left(\varphi_{1}+\varphi_{3}\right)\left(\frac{a_{1}}{b_{1}}\right)^{2^{i}}+\left(\varphi_{2}+\varphi_{3}\right) \frac{a_{1}}{b_{1}}+\varphi_{3}+\varphi_{4}=0
$$

namely,

$$
\left(\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \cdot \frac{a_{1}}{b_{1}}\right)^{2^{i}}+\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \cdot \frac{a_{1}}{b_{1}}=\frac{\varphi_{3}+\varphi_{4}}{\varphi_{3}}
$$

and thus $a_{1}=\alpha b_{1}$ or $\left(\alpha+\frac{\varphi_{3}}{\varphi_{2}+\varphi_{3}}\right) b_{1}$ due to Lemma 14. Therefore, $\operatorname{Det}\left(A_{1}\right)=0$ if and only if $b_{1}=0$ or $a_{1}=\alpha b_{1}$ or $\left(\alpha+\frac{\varphi_{3}}{\varphi_{2}+\varphi_{3}}\right) b_{1}$. Similarly, $\operatorname{Det}\left(B_{1}\right)=0$ if and only if $a_{1}=0$ or $b_{1}=\alpha a_{1}$ or $\left(\alpha+\frac{\varphi_{3}}{\varphi_{2}+\varphi_{3}}\right) a_{1}$.

It is easy to verify that $\operatorname{Det}\left(A_{1}\right)=0$ and $\operatorname{Det}\left(B_{1}\right)=0$ holds at the same time if and only if
(i) $\alpha=1, a_{1}=b_{1}$;
(ii) $\alpha+\frac{\varphi_{3}}{\varphi_{2}+\varphi_{3}}=1, a_{1}=b_{1}$;
(iii) $\alpha\left(\alpha+\frac{\varphi_{3}}{\varphi_{2}+\varphi_{3}}\right)=1, a_{1}=\alpha b_{1}$.

If $\alpha+\frac{\varphi_{3}}{\varphi_{2}+\varphi_{3}}=1$, then $\frac{\varphi_{2}}{\varphi_{3}}=\frac{\alpha}{\alpha+1}$. Recall that (21) holds, namely,

$$
\left(\frac{\varphi_{2}}{\varphi_{3}}\right)^{2}=\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha+\left(\frac{\varphi_{2}+\varphi_{3}}{\varphi_{3}} \alpha\right)^{2}
$$

Plugging $\frac{\varphi_{2}}{\varphi_{3}}=\frac{\alpha}{\alpha+1}$ into the above equation and simplifying, we obtain $\alpha=1$, implying $\frac{\varphi_{3}}{\varphi_{2}+\varphi_{3}}=0$, which is impossible. If $\alpha\left(\alpha+\frac{\varphi_{3}}{\varphi_{2}+\varphi_{3}}\right)=1$, then $\frac{\varphi_{2}}{\varphi_{3}}=\frac{\alpha^{2}+\alpha+1}{\alpha^{2}+1}=\frac{1}{\alpha+1}+\frac{1}{\alpha^{2}+1}+1$, which is also impossible since $\operatorname{Tr}_{q}\left(\frac{\varphi_{2}}{\varphi_{3}}\right)=0$. Therefore, $\operatorname{Det}\left(A_{1}\right)=0$ and $\operatorname{Det}\left(B_{1}\right)=0$ holds at the same time if and only if $\alpha=1, a_{1}=b_{1}$, under which it is clear that $\operatorname{Im}_{V_{i},\left(a_{1}, b_{1}\right)}=\operatorname{Im}_{V_{i},\left(a_{2}, b_{2}\right)}$.

Next, we consider the following two cases:
(I) $\operatorname{Det}\left(B_{1}\right) \neq 0$;
(II) $\operatorname{Det}\left(A_{1}\right) \neq 0$.

It is clear that $\operatorname{Im}_{V_{i},\left(a_{1}, b_{1}\right)}=\operatorname{Im}_{V_{i},\left(a_{2}, b_{2}\right)}$ if there exists some invertible matrix $P$ such that $P A_{1}=A_{2}$ and $P B_{1}=B_{2}$.

As for (i), it suffices to show that

$$
\begin{equation*}
B_{2} B_{1}^{-1} A_{1}=A_{2} \tag{44}
\end{equation*}
$$

After computing, we know that (44) is

$$
\begin{aligned}
& {\left[\begin{array}{l}
b_{21} b_{14} a_{11}+b_{21} b_{12} a_{13}+b_{22} b_{13} a_{11}+b_{22} b_{11} a_{13}, b_{21} b_{14} a_{12}+b_{21} b_{12} a_{14}+b_{22} b_{13} a_{12}+b_{22} b_{11} a_{14} \\
b_{23} b_{14} a_{11}+b_{23} b_{12} a_{13}+b_{24} b_{13} a_{11}+b_{24} b_{11} a_{13}, b_{23} b_{14} a_{12}+b_{23} b_{12} a_{14}+b_{24} b_{13} a_{12}+b_{24} b_{11} a_{14}
\end{array}\right]} \\
& =\operatorname{Det}\left(B_{1}\right)\left[\begin{array}{l}
a_{21}, a_{22} \\
a_{23}, a_{24}
\end{array}\right] .
\end{aligned}
$$

By plugging the expression of $B_{1}$ into the above equation and simplifying, we get

$$
\left\{\begin{array}{l}
b_{14} a_{11}+b_{12} a_{13}=\left(\varphi_{1}+\varphi_{3}\right) a_{1}^{2^{i}+1}+\varphi_{4} a_{1} b_{1}^{2^{i}}+\left(\varphi_{2}+\varphi_{3}\right) b_{1}^{2^{i}+1} \\
b_{14} a_{12}+b_{12} a_{14}=\left(\varphi_{1}+\varphi_{3}\right) a_{1}^{2^{i+1}}+\varphi_{3} a_{1}^{2^{i}} b_{1}^{2^{i}}+\left(\varphi_{1}+\varphi_{3}\right) b_{1}^{2^{i+1}} \\
b_{13} a_{11}+b_{11} a_{13}=\left(\varphi_{2}+\varphi_{3}\right) a_{1}^{2}+\varphi_{3} a_{1} b_{1}+\left(\varphi_{2}+\varphi_{3}\right) b_{1}^{2} \\
b_{13} a_{12}+b_{11} a_{14}=\left(\varphi_{2}+\varphi_{3}\right) a_{1}^{2^{i}+1}+\varphi_{4} a_{1}^{2^{i}} b_{1}+\left(\varphi_{1}+\varphi_{3}\right) b_{1}^{2^{i}+1}
\end{array}\right.
$$

Moreover, we have
1.

$$
\begin{aligned}
& b_{21} b_{14} a_{11}+b_{21} b_{12} a_{13}+b_{22} b_{13} a_{11}+b_{22} b_{11} a_{13} \\
= & \left(\varphi_{3}+\varphi_{4}\right) a_{1}^{2^{i}+2}+\left(\varphi_{4} \alpha^{2^{i}}+\frac{\varphi_{4}\left(\varphi_{2}+\varphi_{3}\right)}{\varphi_{3}} \beta+\frac{\left(\varphi_{2}+\varphi_{3}\right)^{2}\left(\varphi_{1}+\varphi_{3}\right)}{\varphi_{3}}\right) a_{1}^{2} b_{1}^{b_{1}^{i}} \\
+ & \left(\left(\varphi_{2}+\varphi_{3}\right) \alpha^{2^{i}}+\frac{\left(\varphi_{2}+\varphi_{3}\right)^{2}}{\varphi_{3}} \beta+\frac{\left(\varphi_{3}+\varphi_{4}\right)\left(\varphi_{2}+\varphi_{3}\right)\left(\varphi_{1}+\varphi_{3}\right)}{\varphi_{3}}\right) a_{1} b_{1}^{2^{i}+1} \\
+ & \left(\varphi_{3} \alpha+\left(\varphi_{1}+\varphi_{3}\right) \beta+\frac{\left(\varphi_{2}+\varphi_{3}\right)\left(\varphi_{1}+\varphi_{3}\right)^{2}}{\varphi_{3}}\right) a_{1}^{2^{i}+1} b_{1}+\left(\left(\varphi_{2}+\varphi_{3}\right) \alpha+\frac{\left(\varphi_{2}+\varphi_{3}\right)\left(\varphi_{1}+\varphi_{3}\right)}{\varphi_{3}} \beta\right) a_{1}^{2^{i}} b_{1}^{2} \\
= & \left(\varphi_{3}+\varphi_{4}\right) a_{1}^{2^{i}+2}+\left(\varphi_{1}+\varphi_{3}\right) a_{1}^{2} b_{1}^{b_{1}^{i}}+\frac{\left(\alpha^{2}+1\right)\left(\varphi_{2}+\varphi_{3}\right)\left(\varphi_{1}+\varphi_{3}\right)}{\varphi_{3}} a_{1} b_{1}^{i^{i}+1} \\
+ & \frac{\left(\alpha^{4}+\beta^{2}+1\right)\left(\varphi_{2}+\varphi_{3}\right)}{\varphi_{3}} a_{1}^{2^{i}+1} b_{1}+\frac{\left(\alpha^{2}+1\right)\left(\varphi_{2}+\varphi_{3}\right)^{2}}{\varphi_{3}} a_{1}^{2^{i} b_{1}^{2}},
\end{aligned}
$$

2. 

$$
\begin{aligned}
& b_{23} b_{14} a_{11}+b_{23} b_{12} a_{13}+b_{24} b_{13} a_{11}+b_{24} b_{11} a_{13} \\
= & \frac{\left(\varphi_{1}+\varphi_{2}\right)^{2} \beta}{\varphi_{3}} a_{1}^{2^{i}+2}+\frac{\varphi_{4}\left(\varphi_{1}+\varphi_{3}\right) \beta}{\varphi_{3}} a_{1}^{2} b_{1}^{2^{i}}+\frac{\left(\varphi_{1}+\varphi_{3}\right)\left(\varphi_{2}+\varphi_{3}\right) \beta}{\varphi_{3}} a_{1} b_{1}^{2^{i}+1} \\
+ & \left(\varphi_{2}+\varphi_{3}\right) \beta a_{1}^{2^{i}+1} b_{1}+\frac{\left(\varphi_{2}+\varphi_{3}\right)^{2} \beta}{\varphi_{3}} a_{1}^{2^{i}} b_{1}^{2}
\end{aligned}
$$

3. 

$$
\begin{aligned}
& b_{23} b_{14} a_{12}+b_{23} b_{12} a_{14}+b_{24} b_{13} a_{12}+b_{24} b_{11} a_{14} \\
= & \frac{\left(\varphi_{1}+\varphi_{2}\right)^{2} \beta}{\varphi_{3}} a_{1}^{2^{i+1}+1}+\left(\varphi_{1}+\varphi_{3}\right) \beta a_{1}^{2^{i}+1} b_{1}^{2^{i}}+\frac{\left(\varphi_{1}+\varphi_{3}\right)^{2} \beta}{\varphi_{3}} a_{1} b_{1}^{2^{i+1}} \\
+ & \frac{\varphi_{4}\left(\varphi_{2}+\varphi_{3}\right) \beta}{\varphi_{3}} a_{1}^{2^{i+1}} b_{1}+\frac{\left(\varphi_{1}+\varphi_{3}\right)\left(\varphi_{2}+\varphi_{3}\right) \beta}{\varphi_{3}} a_{1}^{2^{i}} b_{1}^{2^{i}+1} .
\end{aligned}
$$

4. 

$$
\begin{aligned}
& b_{21} b_{14} a_{12}+b_{21} b_{12} a_{14}+b_{22} b_{13} a_{12}+b_{22} b_{11} a_{14} \\
= & \left(\varphi_{3}+\varphi_{4}\right) a_{1}^{2^{i+1}+1}+\left(\varphi_{3} \alpha^{2^{i}}+\left(\varphi_{2}+\varphi_{3}\right) \beta+\frac{\left(\varphi_{2}+\varphi_{3}\right)^{2}\left(\varphi_{1}+\varphi_{3}\right)}{\varphi_{3}}\right) a_{1}^{2^{i}+1} b_{1}^{2^{i}} \\
+ & \left(\left(\varphi_{1}+\varphi_{3}\right) \alpha^{2^{i}}+\frac{\left(\varphi_{2}+\varphi_{3}\right)\left(\varphi_{1}+\varphi_{3}\right)}{\varphi_{3}} \beta\right) a_{1} b_{1}^{2^{i+1}} \\
+ & \left(\frac{\left(\varphi_{2}+\varphi_{3}\right)\left(\varphi_{1}+\varphi_{3}\right)^{2}}{\varphi_{3}}+\varphi_{4} \alpha+\frac{\varphi_{4}\left(\varphi_{1}+\varphi_{3}\right)}{\varphi_{3}} \beta\right) a_{1}^{2^{i+1}} b_{1} \\
+ & \left(\left(\varphi_{2}+\varphi_{3}\right)\left(\varphi_{1}+\varphi_{3}\right)+\left(\varphi_{1}+\varphi_{3}\right) \alpha+\frac{\left(\varphi_{1}+\varphi_{3}\right)^{2}}{\varphi_{3}} \beta+\frac{\varphi_{4}\left(\varphi_{2}+\varphi_{3}\right)\left(\varphi_{1}+\varphi_{3}\right)}{\varphi_{3}}\right) a_{1}^{2^{i}} b_{1}^{2^{i}+1} \\
= & \left(\varphi_{3}+\varphi_{4}\right) a_{1}^{2^{i+1}+1}+\frac{\left(\alpha^{2^{i+2}}+\beta^{2}+1\right)\left(\varphi_{1}+\varphi_{3}\right)}{\varphi_{3}} a_{1}^{2^{i}+1} b_{1}^{2^{i}}+\frac{\left(\alpha^{2^{i+1}}+1\right)\left(\varphi_{1}+\varphi_{3}\right)^{2}}{\varphi_{3}} a_{1} b_{1}^{2^{i+1}} \\
+ & \left(\varphi_{2}+\varphi_{3}\right) a_{1}^{2^{i+1}} b_{1}+\frac{\left(\alpha^{2^{i+1}}+1\right)\left(\varphi_{1}+\varphi_{3}\right)\left(\varphi_{2}+\varphi_{3}\right)}{\varphi_{1} b_{1}^{2^{i+1}},}
\end{aligned}
$$

Furthermore, after computing and simplifying, we have
1.

$$
\begin{aligned}
& \operatorname{Det}\left(B_{1}\right) a_{21} \\
= & \left(\varphi_{3}+\varphi_{4}\right) a_{1}^{2^{i}+2}+\left(\varphi_{1}+\varphi_{3}\right) a_{1}^{2} b_{1}^{2^{i}}+\frac{\left(\alpha^{2}+1\right)\left(\varphi_{2}+\varphi_{3}\right)\left(\varphi_{1}+\varphi_{3}\right)}{\varphi_{3}} a_{1} b_{1}^{2^{i}+1} \\
+ & \frac{\left(\alpha^{4}+\beta^{2}+1\right)\left(\varphi_{2}+\varphi_{3}\right)}{\varphi_{3}} a_{1}^{2^{i}+1} b_{1}+\frac{\left(\alpha^{2}+1\right)\left(\varphi_{2}+\varphi_{3}\right)^{2}}{\varphi_{3}} a_{1}^{2^{i}} b_{1}^{2}
\end{aligned}
$$

2. 

$$
\begin{aligned}
& \operatorname{Det}\left(B_{1}\right) a_{22} \\
= & \left(\varphi_{3}+\varphi_{4}\right) a_{1}^{2^{i+1}+1}+\frac{\left(\alpha^{2^{i+2}}+\beta^{2}+1\right)\left(\varphi_{1}+\varphi_{3}\right)}{\varphi_{3}} a_{1}^{2^{i}+1} b_{1}^{2^{i}}+\frac{\left(\alpha^{2^{i+1}}+1\right)\left(\varphi_{1}+\varphi_{3}\right)^{2}}{\varphi_{3}} a_{1} b_{1}^{2^{i+1}} \\
+ & \left(\varphi_{2}+\varphi_{3}\right) a_{1}^{2^{i+1}} b_{1}+\frac{\left(\alpha^{2^{i+1}}+1\right)\left(\varphi_{1}+\varphi_{3}\right)\left(\varphi_{2}+\varphi_{3}\right)}{\varphi_{3}} a_{1} b_{1}^{2^{i+1}}
\end{aligned}
$$

3. 

$$
\begin{aligned}
& \operatorname{Det}\left(B_{1}\right) a_{23} \\
= & \frac{\left(\varphi_{1}+\varphi_{2}\right)^{2} \beta}{\varphi_{3}} a_{1}^{2^{i}+2}+\frac{\varphi_{4}\left(\varphi_{1}+\varphi_{3}\right) \beta}{\varphi_{3}} a_{1}^{2} b_{1}^{2^{i}}+\frac{\left(\varphi_{1}+\varphi_{3}\right)\left(\varphi_{2}+\varphi_{3}\right) \beta}{\varphi_{3}} a_{1} b_{1}^{2^{i}+1} \\
+ & \left(\varphi_{2}+\varphi_{3}\right) \beta a_{1}^{2^{i}+1} b_{1}+\frac{\left(\varphi_{2}+\varphi_{3}\right)^{2} \beta}{\varphi_{3}} a_{1}^{2^{i}} b_{1}^{2}
\end{aligned}
$$

4. 

$$
\begin{aligned}
& \operatorname{Det}\left(B_{1}\right) a_{24} \\
= & \frac{\left(\varphi_{1}+\varphi_{2}\right)^{2} \beta}{\varphi_{3}} a_{1}^{2^{i+1}+1}+\left(\varphi_{1}+\varphi_{3}\right) \beta a_{1}^{2^{i}+1} b_{1}^{2^{i}}+\frac{\left(\varphi_{1}+\varphi_{3}\right)^{2} \beta}{\varphi_{3}} a_{1} b_{1}^{i^{i+1}} \\
+ & \frac{\varphi_{4}\left(\varphi_{2}+\varphi_{3}\right) \beta}{\varphi_{3}} a_{1}^{2^{i+1}} b_{1}+\frac{\left(\varphi_{1}+\varphi_{3}\right)\left(\varphi_{2}+\varphi_{3}\right) \beta}{\varphi_{3}} a_{1}^{2^{i}} b_{1}^{2^{i}+1} .
\end{aligned}
$$

Hence, it follows that

$$
\begin{aligned}
& {\left[\begin{array}{l}
b_{21} b_{14} a_{11}+b_{21} b_{12} a_{13}+b_{22} b_{13} a_{11}+b_{22} b_{11} a_{13}, b_{21} b_{14} a_{12}+b_{21} b_{12} a_{14}+b_{22} b_{13} a_{12}+b_{22} b_{11} a_{14} \\
b_{23} b_{14} a_{11}+b_{23} b_{12} a_{13}+b_{24} b_{13} a_{11}+b_{24} b_{11} a_{13}, b_{23} b_{14} a_{12}+b_{23} b_{12} a_{14}+b_{24} b_{13} a_{12}+b_{24} b_{11} a_{14}
\end{array}\right]} \\
& =\operatorname{Det}\left(B_{1}\right)\left[\begin{array}{l}
a_{21}, a_{22} \\
a_{23}, a_{24}
\end{array}\right] .
\end{aligned}
$$

and Eq. (44) holds.
As for (ii), we need to show that

$$
\begin{equation*}
A_{2} A_{1}^{-1} B_{1}=B_{2} \tag{45}
\end{equation*}
$$

whose proof can be obtained through just changing $a_{1}$ and $b_{1}$ in the proof of (44).
Therefore, for any $\left(a_{1}, b_{1}\right),\left(a_{2}, b_{2}\right) \in \mathbb{F}_{q}^{2} \backslash\{(0,0)\}$ satisfying $B_{V_{i},\left(a_{1}, b_{1}\right)}\left(a_{2}, b_{2}\right)=(0,0), \operatorname{Im}_{V_{i},\left(a_{1}, b_{1}\right)}=$ $\operatorname{Im}_{V_{i},\left(a_{2}, b_{2}\right)}$ holds and by Lemma 5 , we know that the boomerang uniformity of $V_{i}$ is 4 .

Remark 17 From [13], we are aware that the authors are independently working on the same problem as in this paper. Their techniques in the proof are different from ours in the early version [11] of this paper.

Remark 18 It's worth pointing out that from the experimental results by MAGMA, the set $\Gamma$ in Theorem 1 covers all the coefficients $\alpha, \beta \in \mathbb{F}_{q}^{*}$ that yield permutations $V_{i}(x, y)$ with boomerang uniformity 4 with $q=2^{3}, 2^{5}$. We therefore propose the following conjecture and invite interested readers to attack it.

Conjecture 19 Let $q=2^{n}$ with $n$ odd, $\operatorname{gcd}(i, n)=1$ and $V_{i}:=\left(R_{i}(x, y), R_{i}(y, x)\right)$ with $R_{i}(x, y)=$ $(x+\alpha y)^{2^{i}+1}+\beta y^{2^{i}+1}$. If $V_{i}$ is a permutation over $\mathbb{F}_{q}^{2}$ with boomerang uniformity 4 , then the coefficients $\alpha, \beta$ are taken from the set $\Gamma$ defined as in (1).

## 5 Conclusions

In this paper, we apply the butterfly structure by Canteaut et. al in constructing cryptographically strong permuations. The open butterfly does not seem to generate permutations with boomerang uniformity 4 according to numerical result. Based on an intensive study on the coefficients of $R_{i}(x, y)=$ $(x+\alpha y)^{2^{i}+1}+\beta y^{2^{i}+1}, \operatorname{gcd}(i, n)=1$, over $\mathbb{F}_{2^{n}}$, we provide a sufficient condition on $\alpha, \beta$ such that $V_{i}(x, y)=\left(R_{i}(x, y), R_{i}(y, x)\right)$ is a permutation over $\mathbb{F}_{2^{2 n}}$ with boomerang uniformity 4 . The proposed condition seems to be necessary according to numeric results and a conjecture on the observation is given.

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