

Copyright

by

Lu Gao

2011

**The Dissertation Committee for Lu Gao Certifies that this is the approved  
version of the following dissertation:**

**Optimal Infrastructure Maintenance Scheduling Problem Under  
Budget Uncertainty**

**Committee:**

---

Zhanmin Zhang, Supervisor

---

Carlos Caldas

---

Stephen Donald

---

Randy B. Machehmel

---

Michael R. Murphy

---

C. Michael Walton

# **Optimal Infrastructure Maintenance Scheduling Problem under Budget Uncertainty**

by

Lu Gao B.E.; M.S.E.

**Dissertation**

Presented to

The Faculty of the Graduate School of

The University of Texas at Austin

in Partial Fulfillment

of the Requirements

for the Degree of

**Doctor of Philosophy**

**The University of Texas at Austin**

**August 2011**

## **Dedication**

This dissertation is dedicated to my parents and wife for their unconditional love and support.

## **Acknowledgements**

I am heartily thankful to my supervisor, Dr. Zhanmin Zhang, whose encouragement, supervision and support from the preliminary to the concluding level enabled me to develop an understanding of the subject. I also want to thank my dissertation committee members Dr. Carlos Caldas, Dr. Stephen Donald, Dr. Randy B. Machemehl, Dr. Michael R. Murphy and Dr. C. Michael Walton for their advice and support.

I wish to thank my friends, Cai Shuwen, Chen Zhu, Dong Changgui, Guo Runhua, Hong Feng, José Pablo Aguiar-Moya, Liang Hangwen, Li Qi, Luo Wenjie, Yu Ke, Yu Yao, Xie Chi, Wang Pei, Wu Hui, and Zhang Guohui, for helping me get through the difficult times, and for all the emotional support, comraderie, entertainment, and caring they provided.

Finally, I would like to thank my parents for their unconditional love and support. I would also like to express my deep appreciation and gratitude to my wife, Yang Yufan, for her compassionate love and encouragement.

# **Optimal Infrastructure Maintenance Scheduling Problem under Budget Uncertainty**

Lu Gao, Ph.D.

The University of Texas at Austin, 2011

Supervisor: Zhanmin Zhang

This research addresses the infrastructure maintenance scheduling problems under budget uncertainty. Infrastructure agencies usually face budget uncertainties that will eventually lead to suboptimal planning if maintenance decisions are made without taking the uncertainty into consideration. It is important for decision makers to adopt maintenance scheduling policies that take future budget uncertainty into consideration.

The author proposes a multistage, stochastic linear programming model to address this problem. The author also develops solution procedures using the augmented Lagrangian decomposition algorithm and scenario reduction method. A case study exploring the computational characteristics of the proposed methods is conducted and the benefit of using the stochastic programming approach is discussed. In the case study, the road network in Dallas District is used with data taken from the Texas Department of Transportation's Pavement Management Information System. The case study results reveal that the

stochastic programming solutions tend to allocate more resources to preventive maintenance than deterministic solutions that ignore the uncertainty information. The proposed methodology can help decision makers effectively obtain optimal maintenance plan under budget uncertainty.

# TABLE OF CONTENTS

<b>TABLE OF CONTENTS .....</b>	<b>viii</b>
<b>LIST OF TABLES.....</b>	<b>xi</b>
<b>LIST OF FIGURES .....</b>	<b>xii</b>
<b>CHAPTER 1 INTRODUCTION .....</b>	<b>1</b>
1.1 Infrastructure Asset Management.....	1
1.1.1 Data Collection .....	3
1.1.2 Performance Modeling .....	4
1.1.3 Program Optimization.....	4
1.1.4 Feedback.....	5
1.2 Infrastructure Asset Management Maintenance Scheduling Problem ....	5
1.2.1 Project-Level Problem .....	6
1.2.2 Network-Level Problem .....	7
1.3 Motivation of this Research .....	8
1.4 Objective.....	10
1.5 Dissertation Organization .....	10
<b>CHAPTER 2 LITERATURE REVIEW OF INFRASTRUCTURE</b>	
<b>PERFORMANCE MODELS AND MAINTENANCE SCHEDULING MODELS ..</b>	<b>12</b>
2.1 Performance Models .....	12
2.1.1 Markov Chain Model.....	14
2.1.2 Reliability Model (or Survival Model) .....	17



2.1.3	Discrete Choice Model .....	19
2.1.4	Deterministic Models .....	21
2.2	Maintenance Scheduling Models.....	22
2.2.1	Markov Chain Based Linear Programming (LP) .....	23
2.2.2	Integer Programming (IP) .....	24
2.2.3	Reliability Model .....	29
2.2.4	Optimal Control Model.....	30
2.3	Infrastructure Maintenance Scheduling Models Considering Budget Uncertainty .....	34
<b>CHAPTER 3 PROBLEM FORMULATION .....</b>		<b>37</b>
3.1	Notations .....	37
3.2	Deterioration Modeling.....	39
3.3	Maintenance Planning Model with Deterministic Budgets .....	41
3.4	Maintenance Planning Model Under Budget Uncertainty .....	43
<b>CHAPTER 4 SOLUTION PROCEDURE .....</b>		<b>48</b>
4.1	Augmented Lagrangian Decomposition (ALD) .....	49
4.2	Application of ALD to Stochastic Programming .....	52
4.3	Scenario Reduction (SR) .....	54
<b>CHAPTER 5 CASE STUDY.....</b>		<b>56</b>
5.1	Case Study Data Set .....	56
5.1.1	Size of the Network .....	57
5.1.2	Planning Horizon .....	57
5.1.3	Performance Indicator .....	58

5.1.4	Transition Probability .....	60
5.1.5	Maintenance Effect.....	62
5.1.6	Maintenance Cost (Agency Cost) .....	64
5.2	Solution of ALD.....	65
5.3	Deterministic Solution (EV).....	67
5.4	Computational Comparison (ALD, EV and SR) .....	71
<b>CHAPTER 6 CONCLUSIONS AND RECOMMENDATIONS .....</b>		<b>75</b>
6.1	Summary .....	75
6.2	Conclusions.....	75
6.3	Recommendations for Future Research .....	76
6.3.1	Stochastic Integer Programming .....	76
6.3.2	Uncertainties other than Budget .....	77
6.3.3	Different Ownership.....	78
6.3.4	Balance between different regions .....	78
6.3.5	Multiple objectives .....	79
<b>REFERENCES .....</b>		<b>80</b>
<b>VITA.....</b>		<b>89</b>

## LIST OF TABLES

Table 3.1 Notation .....	37
Table 5.1 Road Network Length .....	57
Table 5.2 PMIS Condition Scores.....	58
Table 5.3 Road Network Initial Condition (%) .....	59
Table 5.4 Road Condition Requirements .....	60
Table 5.5 TPM for Road Group I.....	60
Table 5.6 TPM for Road Group II.....	61
Table 5.7 TPM for Road Group III.....	61
Table 5.8 Maintenance Treatments Effect .....	63
Table 5.9 Maintenance Treatment Costs .....	64
Table 5.10 Computational Characteristics .....	65
Table 5.11 Maintenance Plan of Deterministic Solution.....	69
Table 5.12 Maintenance Plan Comparison of Year 1 between EV and SP .	71
Table 5.13 Computational Characteristics Comparison .....	74

## LIST OF FIGURES

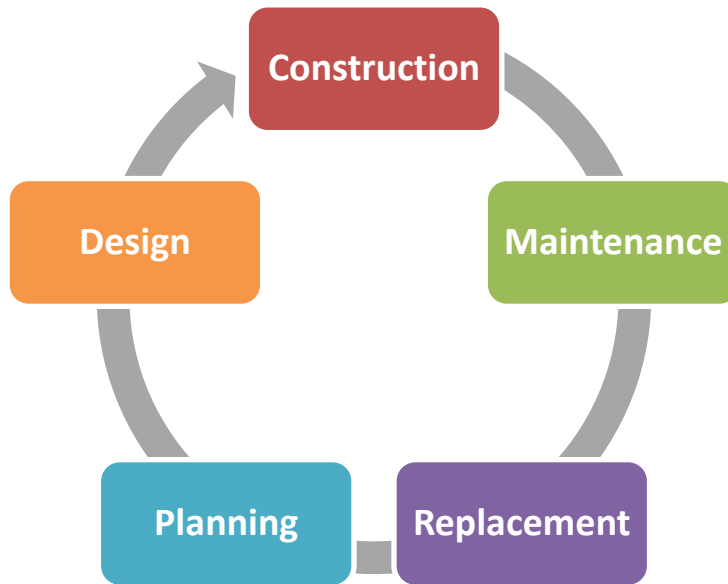
Figure 1-1 Life Cycle Phases of Infrastructure Asset Management .....	2
Figure 1-2 Basic Elements of Infrastructure Asset Management .....	3
Figure 1-3 Project-Level Maintenance Scheduling Problem .....	7
Figure 2-1 Markov Chain Performance Model .....	16
Figure 2-2 Reliability Curve of an Infrastructure Facility .....	19
Figure 2-3 Infrastructure Condition Trend Trajectory .....	32
Figure 3-1 Scenario Tree .....	45
Figure 4-1 Scenario Tree after Decomposition .....	52
Figure 5-1 Number of Jacobi Steps in Each Outer Loop .....	66
Figure 5-2 Number of Violated Nonanticipativity Constraints .....	67
Figure 5-3 Scenario Tree of EV Approach .....	68
Figure 5-4 Scenario Tree before Reduction .....	72
Figure 5-5 Scenario Tree after Reduction .....	73

# CHAPTER 1 INTRODUCTION

## 1.1 INFRASTRUCTURE ASSET MANAGEMENT

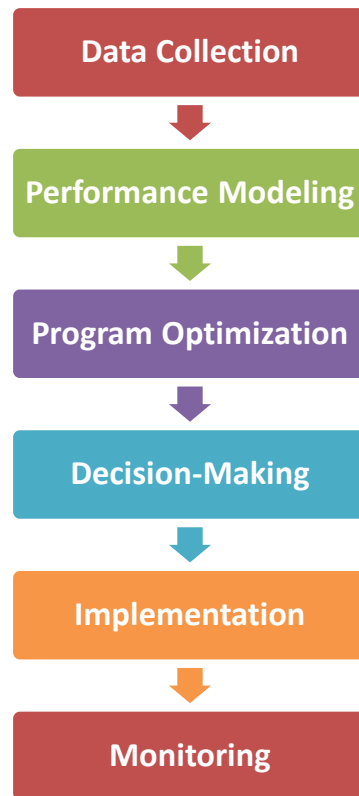
Infrastructure asset management (IAM) is a systematic approach to maintaining, upgrading, and operating infrastructure facilities cost-effectively. Examples of infrastructure assets include pavements, bridges, drainage culverts, storm drainage systems, traffic signals, traffic signs, traffic striping, ITS infrastructure, safety rest areas and roadside. IAM combines together engineering principles, sound business practices and economic theory. The goal of infrastructure asset management is to effectively manage large and complex infrastructure systems in an integrated manner.

Generally, the management process of IAM focuses on the stages of maintenance, rehabilitation, and replacement. Mathematical models and computer software tools are employed to help decision makers preserve and extend the service life of infrastructure facilities. However, in the broadest sense, infrastructure management covers all phases of an infrastructure project from planning, design, construction, maintenance to disposal as illustrated in Figure 1-1.



**Figure 1-1 Life Cycle Phases of Infrastructure Asset Management**

All infrastructure facilities deteriorate over time due to various reasons. As infrastructure facilities deteriorate, the cost to operate and maintain them increases. Therefore, managing maintenance activities for large-scale infrastructure systems is a difficult task. Many projects and interests compete for the limited resources allocated to different programs. Many factors are involved in the decision-making process of infrastructure asset management. The basic elements of infrastructure asset management are shown in Figure 1-2.



**Figure 1-2 Basic Elements of Infrastructure Asset Management**

### **1.1.1 Data Collection**

Data collection provides the decision makers with the information about the condition of the infrastructure system. Moreover, the data collected will also provide basic information about the location and inter-connectivity of each management segment. The minimum data required for each management segment generally include identification, location, size, importance such as functional classification, material type, usage levels and date of construction or last major repair. For example, for a network of pavements, the data needed include various types of distress (such as rutting, roughness, slab faulting),

pavement strength, and so on. Such data will be used in preparing the needs analysis and maintenance activity scheduling. Moreover, data about construction quality and maintenance history can also be inventoried.

### **1.1.2 Performance Modeling**

Infrastructure deterioration is a complicated, dynamic, and stochastic process affected by various factors such as design, climate conditions (e.g., rainfall, temperature and amount of sunlight), material, structural capacities, and some unobserved factors. Accurate prediction of infrastructure performance is critical to infrastructure asset management agencies. Reliable and accurate prediction of infrastructure performance can save significant amounts of money by helping plan maintenance and rehabilitation activities. Performance models (or deterioration models) can be developed by using historical data. Usually, data points affected by maintenance activities are excluded when used in the development of performance models, in order to obtain the true deterioration process of the facility.

### **1.1.3 Program Optimization**

By using the performance model, the condition of infrastructure facilities can be projected into the future. However, the projected condition may not satisfy the decision makers' requirement. Therefore, in order to achieve an established goal, the maintenance plan needs to be adjusted and facilities will be selected to receive maintenance and rehabilitation treatments during the planning horizon.



Once the decision maker has determined the funding needed to maintain the system in desired condition, the identified funding needs will be compared to the funding available. If the available funding is less than what is needed for any of the years in the analysis period, optimization technique (or other methods) should be used. The goal of using optimization technique is to find the best funding allocation strategy that will provide the greatest overall return.

#### **1.1.4 Feedback**

Feedback refers to the transfer of part of the output to the input. A feedback system ensures continual feedback of information to monitor infrastructure system conditions. Feedback system also allows the evaluation of other aspects of the infrastructure system, including life-cycle costs and effectiveness of maintenance and rehabilitation treatments. In general, the feedback system provides the information to evaluate the reliability of past estimates.

In the context of a pavement management system, the feedback process allows the actual costs of maintenance to be compared with those obtained in the prediction analysis. Moreover, evaluations from field observations of pavement conditions can also be compared with those predicted by performance models.

## **1.2 INFRASTRUCTURE ASSET MANAGEMENT MAINTENANCE SCHEDULING PROBLEM**

An efficient transportation infrastructure network is vital to a society's economic and social development. Infrastructure maintenance scheduling is one of the

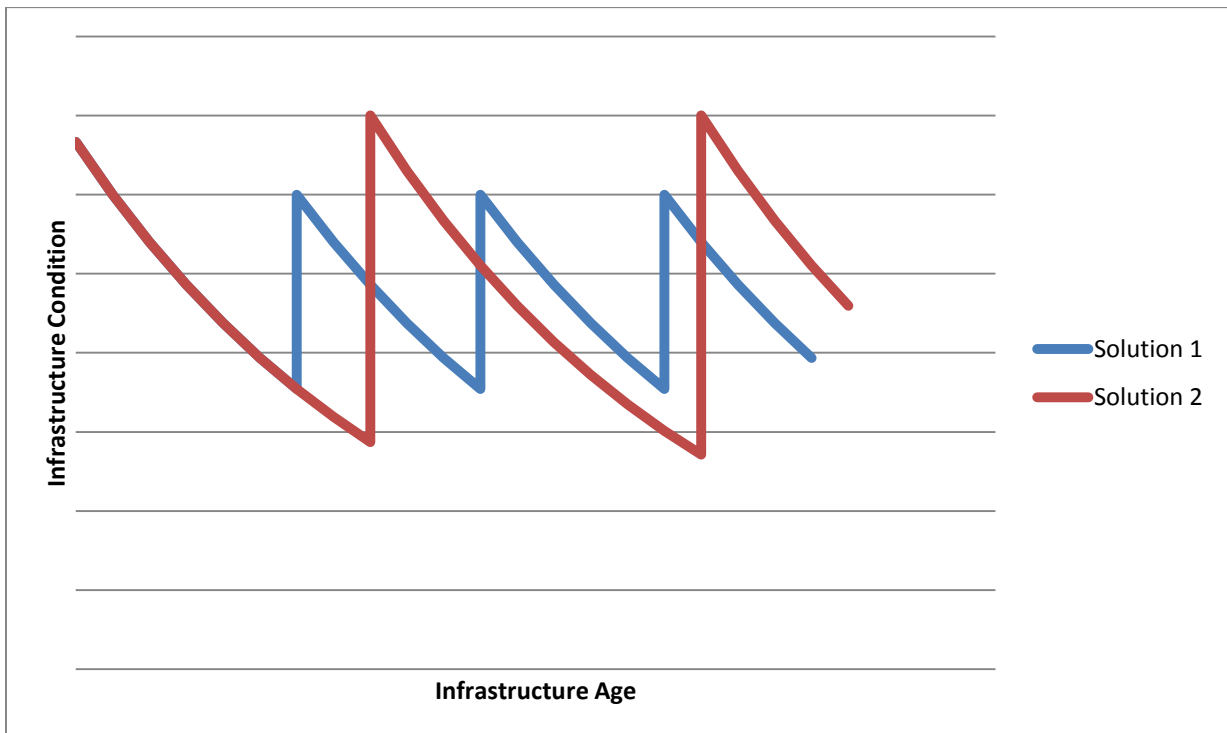
most important components of infrastructure management. It is the process of developing alternative maintenance schedules and determining the best solution to ensure desired level of service. For many infrastructure facilities, the service life can be extended beyond the original design life by applying maintenance treatments. A maintenance schedule is the selection of a sequence of maintenance treatments over the planning period. Life-cycle cost analysis is usually used in choosing between various alternatives. Both user and agency costs should be considered in scheduling maintenance activities.

Maintenance options of an infrastructure facility include routine, preventive, reactive, and other rehabilitation and replacement techniques. Maintenance expenditure is one of the biggest infrastructure investments. From a mathematical point of view, there are two types of maintenance scheduling problems. The first one is the network-level problem, where decision-makers face great challenges of determining which facility should be repaired, when and how repairs should be carried out, and what treatment should be used. The other is the project-level maintenance problem, in which the maintenance scheduling of only one facility is considered.

### **1.2.1 Project-Level Problem**

Project-level problem is the foundation of the network-level problem. Project-level problem is to determine the best maintenance schedule for a single facility by considering both user and agency cost. A typical project-level maintenance scheduling problem is illustrated in Figure 1-3. As showing in the figure, there

are two maintenance scheduling options (blue and red). It requires a life-cycle cost analysis before the decision makers can make a choice between them.



**Figure 1-3 Project-Level Maintenance Scheduling Problem**

### 1.2.2 Network-Level Problem

Network-Level scheduling problems deal with systems of multiple facilities. The mathematical model of a network-level scheduling problem is usually constructed as a combination of models of many project-level problems. Therefore, network-level problems are more complex and more difficult to solve than project-level problems. In a network-level maintenance scheduling problem, decision-makers

usually have to face large-scale systems of facilities. The purpose of the network-level maintenance scheduling problem is normally to identify the fund needs and determine location and timing of maintenance treatments for the whole system.

### **1.3 MOTIVATION OF THIS RESEARCH**

The U.S. population is expected to grow by 100 million during the next 30 years. However, the current investment levels are not keeping pace with the increased usage and deterioration of the nation's infrastructure network. For example, nearly 161,750 miles of federal-aid highways have pavement rated unacceptable and 153,990 bridges nationwide are structurally deficient or functionally obsolete, according to U.S. Department of Transportation (U.S. DOT) data. This situation is not going to change in the next ten years unless steps are taken to improve how available funds are used and to increase the amount of funds to meet system needs. According to data from the 2006 U.S. DOT "Condition & Performance" report, the federal share of highway investment needed just to maintain highway conditions and performance will grow to almost \$62 billion by 2015. Moreover, the U.S. Treasury estimates that the revenues into the Highway Account will grow to just under \$42 billion by 2015. The gap between projected revenues and minimum investment needs average \$19 billion per year.

A recent study conducted by Texas Department of Transportation shows that as a result of use and age, Texas' highway infrastructure is showing signs of deterioration. The total revenue available in Texas for pavement and bridge maintenance plus additional capacity is expected to be \$100 billion from 2011 to

2035. The estimated funding gaps will range from \$74 billion to \$170 billion from 2011 to 2035 (Texas 2030 Committee, 2011). Under these situations, how to effectively use the limited funding on transportation infrastructure is a cutting-edge problem.

A number of mathematical models have been developed for scheduling infrastructure system maintenance activities under budget constraints. Most of the approaches treat the annual budget as a fixed amount. An underlying assumption is that actual funds to support the maintenance activities would never deviate from the original expectation. However, this assumption is often unrealistic because the funding allocated to infrastructure maintenance program is subject to uncertainty due to various financial and political risks. Moreover, the funding for maintenance usually has to compete with other activities, e.g., capacity expansion projects. Consequently, the actual amount of money distributed to the maintenance activities may deviate from the original estimate. Therefore, if the funding falls short for some of the years during the planning period, part of the planned maintenance activities might be forced to postpone, leading to inevitable condition deviation from the expectation. As a result, ignoring the random characteristics of future budget may limit the usefulness of the optimal scheduling solution. It is therefore without doubt that the assumption of deterministic budget is questionable in practice.

#### **1.4 OBJECTIVE**

The objective of this research is to develop a solution framework for network-level infrastructure maintenance scheduling problems under budget uncertainty. The problem was formulated as a multistage, linear stochastic programming model. The proposed model differs from its deterministic counterpart in that it attempts to find the optimal maintenance scheduling plan given the information that future funding is uncertain. The author uses the augmented Lagrangian decomposition method and the scenario reduction method to solve the stochastic programming problem. The usefulness and efficiency of the proposed model will be tested in a road network maintenance scheduling problem.

#### **1.5 DISSERTATION ORGANIZATION**

The organization of this dissertation is as follows: Chapter 1 introduces the problem and research motivation . Chapter 2 focuses on reviewing the literature of infrastructure performance models and maintenance scheduling models. In this chapter, models developed in previous works are classified into different categories. For each category, the advantages and disadvantages are discussed and summarized. Chapter 3 describes the methodologies of formulating the infrastructure maintenance scheduling problem under both deterministic and stochastic conditions. Chapter 4 discusses the solution methods used in this research to solve the multi-stage stochastic programming problem. Chapter 5 presents the application of the models and algorithms developed in Chapter 3

and 4 to the PMIS data set. The optimal solution results are discussed. Chapter 6 summarizes the research effort and presents the conclusions.

## **CHAPTER 2      LITERATURE REVIEW OF INFRASTRUCTURE PERFORMANCE MODELS AND MAINTENANCE SCHEDULING MODELS**

### **2.1 PERFORMANCE MODELS**

In infrastructure asset management, performance models are used to predict future conditions and to help schedule maintenance activities. The effectiveness of maintenance planning in infrastructure management depends on the accuracy of the predicted future condition of infrastructure facilities. If the performance models used in planning maintenance activities cannot accurately represent the actual deterioration process, the planned maintenance activities might not yield the expected results, leading to suboptimal use of resources.

Performance models can be classified into two categories: deterministic and probabilistic. In deterministic models, the future condition of a facility is predicted as an exact value using historical data. In probabilistic models, the future condition of a facility is predicted by estimating the probability with which the facility would change to a particular condition state. Probabilistic models are usually associated with the discretization of the condition indicator. Moreover, probabilistic models can also be used to describe the deterioration process of multiple facilities, i.e., the whole system.

Most of the performance models developed in the early stages of infrastructure asset management research are deterministic (see AASHO (1962), Garcia and Riggins (1984), Paterson (1987), for example). Such models tend to



ignore the stochastic nature of the infrastructure deterioration process. However, infrastructure deterioration is a complex process associated with many uncertainties, which mostly come from three sources. The first source concerns measurement errors, which can cause a high degree of prediction uncertainty (Humplick, 1992). The second source of uncertainty is the inherent randomness of the facility deterioration process. The third source is associated with latent factors (e.g., construction quality), which are difficult to observe and quantify individually. Probabilistic models are developed to help take these uncertainties into consideration when modeling infrastructure deterioration.

A popular example of probabilistic performance models is the one based on the Markov Chain theory, in which the deterioration process is characterized by transition between different condition states. Markov Chain can be used in modeling both single facility (e.g., pavement, bridge) and system (e.g., pavement network). For example, Golabi et al. (1982) proved the effectiveness of using the Markov Chain method by developing Markov Chain performance models in Arizona. The core of the Markov Chain models is the development of the transition probabilities. A number of methods including the expected-value method by Butt et al. (1987) and Jiang et al. (1989) and the proportion method by Wang et al. (1994) have been employed to calculate the transition probabilities. Another way of calculating the transition probabilities is the simulation approach of utilizing the design equations (Gao et al. 2007). More examples of performance models based on Markov Chain theory can be found in Madanat et al. (1995, 1997).

Another type of probabilistic performance model is the reliability model (also called duration model or survival model). For example, Prozzi and Madanat (2000) developed a duration model to predict the number of axle load repetitions a pavement can withstand before its serviceability drops below an acceptable level. Zhang and Damnjanovic (2006) developed a model to predict the reliability of a pavement by using design equations. The limitation of reliability model is that the condition of an infrastructure facility (e.g., a pavement section) is usually characterized by multiple condition states. As a result, using only two states (survival and failure) cannot fully characterize the changing of the facility condition. The reliability model is more suitable for modeling a specific distress failure mode, in which the development of the distress cannot be easily observed until it reaches a certain level (see Wang et al. (2005), for example). It can also be used in such scenarios that the failure of a facility (e.g., a bridge) has significant consequence. In such case, decision makers can better understand the risk if reliability model is used.

In the rest of this section, major existing performance models in the literature are discussed in details.

### **2.1.1 Markov Chain Model**

The Markov Chain model used in infrastructure deterioration modeling is characterized by the following features. First of all, the Markov process is discrete in time. Secondly, the Markov process has a countable state space. Finally, the Markov process satisfies the Markovian property. The Markovian

property is said to be satisfied if the future state of the process depends only on its present state, but not on its past states. This property is satisfied in performance modeling if the future condition of the facility depends only on its present condition and not on its past condition. The Markov Chain model can be used for both project-level (single facility) and network-level (multiple facilities) performance modeling (Gao et al. 2007). For single facility modeling, the solution represents the probabilities of the facility being in different condition states. For network modeling, the solution represents the proportions of the network being in different condition states.

To apply the Markov Chain mode (Figure 2-1), the condition of a facility is first discretized into  $n$  states. Hence, facility condition at different time periods can be represented by a condition state probability vector:

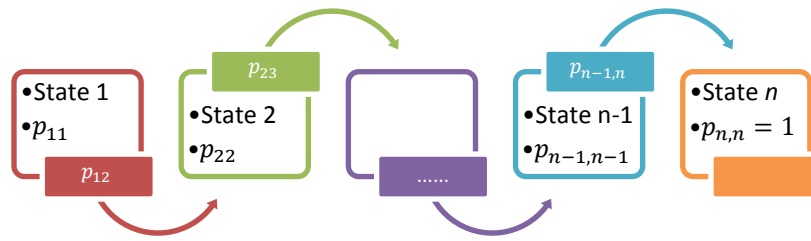
$$C(t) = [c_1(t), c_2(t), \dots, c_n(t)] \quad (2.1)$$

where:

$C(t)$  = condition state probability vector of a facility at time period  $t$ ;

$c_i(t)$  = probability that a facility stays in state  $i$  at time period  $t$ ,  $i = 1, 2, \dots, n$  and

$$\sum_{i=1}^n c_i(t) = 1.$$



**Figure 2-1 Markov Chain Performance Model**

The deterioration process of an infrastructure facility can be expressed as the change of the elements in the condition state probability vector. A transition probability matrix  $P$  is used to calculate this change. In the Markovian process, it is assumed that the future condition states of a facility depend only on its current condition state. Any experience before it has no impact on the future condition. Therefore, in order to calculate future condition state probability, only the present condition state probability vector and  $P$  are needed.

$$P = \begin{pmatrix} p_{11} & p_{12} & 0 & \cdots & 0 & 0 \\ 0 & p_{22} & p_{23} & \cdots & 0 & 0 \\ 0 & 0 & p_{33} & \cdots & \vdots & \vdots \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & \cdots & p_{n-1,n-1} & p_{n-1,n} \\ 0 & 0 & 0 & \cdots & 0 & p_{n,n} = 1 \end{pmatrix} \quad (2.2)$$

where,

$p_{ij}$  = probability that the facility will deteriorate from state  $i$  to state  $j$  at time period  $t$ ;

Because a facility cannot improve to a better condition state by itself, the elements  $p_{ij}$  are replaced by 0 for  $i > j$ . Furthermore, the value of 1 in the last row of the matrix corresponding to state  $n$  indicates that the condition cannot deteriorate further. From all the above, the future condition probability can be calculated by:

$$C(t + \Delta t) = C(t)P^{\Delta t} \quad (2.3)$$

### 2.1.2 Reliability Model (or Survival Model)

Survival analysis is a branch of statistics which deals with the counting of deaths and failures. More generally, survival analysis involves the modeling of time to event. Survival model attempts to answer questions of what is the probability of a subject surviving over a certain time period; and what are the variables that affect the failure. The primary goal in using survival models to analyze infrastructure condition data is to assess the dependence of time-to-failure on external variables. One way to explore the relationship of covariates and time-to-failure is to use a regression model in which the failure time has a probability distribution that depends on the covariates.

The specific feature that distinguishes survival model from classical statistical model is data censoring. Usually, the failure time is unknown for some of the facilities. The only information available is that the facility has survived up to a certain time. This type of censoring is called right censoring. For right-censored data, the actual information of the  $i^{th}$  facility  $i=1, \dots, n$  is contained in

the pair  $(t_i, d_i)$ , where  $t_i$  is the failure time and  $d_i$  is the censoring indicator, taking the value one if the event has been observed (failed), otherwise  $d_i$  takes value zero (censored). Then the censoring indicator can be expressed in Eq.(2.4).

$$d_i = \begin{cases} 1 & \text{if } t_i \leq c_i \\ 0 & \text{if } t_i > c_i \end{cases} \quad (2.4)$$

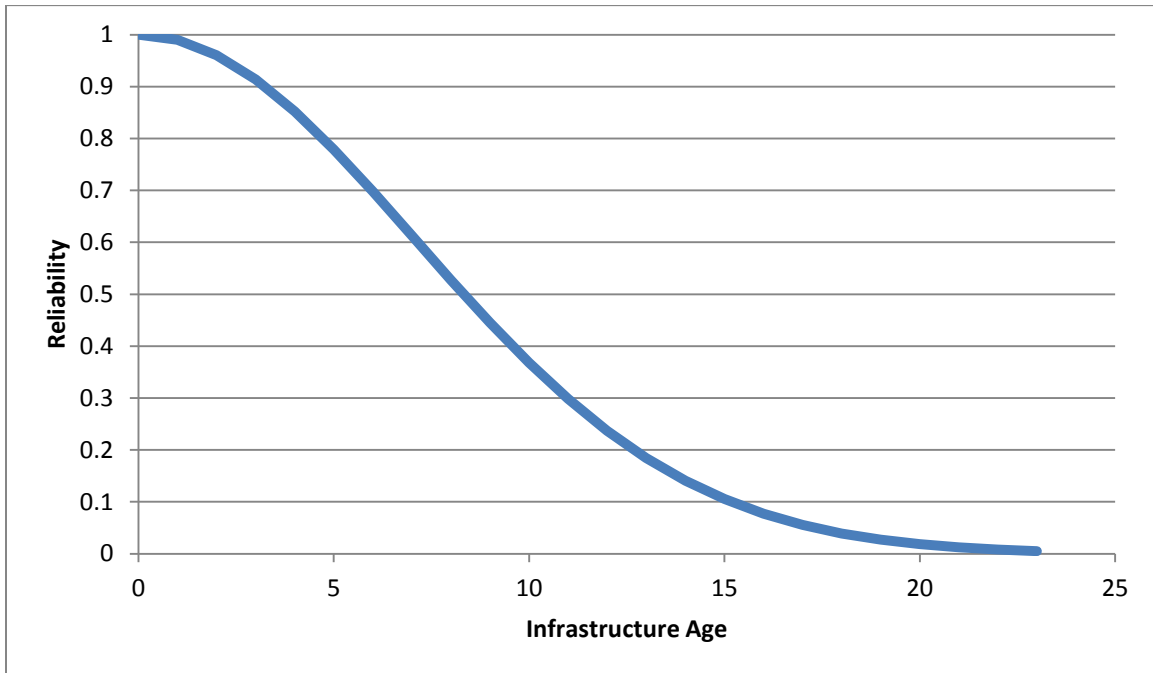
where,  $c_i$  is the censoring time.

For a random time-to-failure,  $T$ , the probability density function of  $T$  is defined as  $f(t)$  and the cumulative distribution function as  $F(t) = P(T \leq t)$ . Two other functions that are useful in this context are the survival function  $S(t) = P(T > t) = 1 - F(t)$ , and the hazard function  $h(t) = f(t)/S(t)$ , which can be interpreted as the instantaneous rate of failure given survival up until time  $t$ .

One of the survival models that have been used in infrastructure performance modeling is the Proportional Hazard (PH) model (Gao et al. 2011). In general, a PH model with covariates can be written as:

$$h(t_i) = h_0(t_i) \exp(\mathbf{x}_i \boldsymbol{\beta}) \quad (2.5)$$

where  $h_0(t)$  is the baseline hazard function representing the deterioration rate of the facility;  $\boldsymbol{\beta}$  is the parameter vector and  $\mathbf{x}_i$  is the covariates vector of the  $i$ th observation. A typical survival curve of an infrastructure facility is shown in Figure 2-2.



**Figure 2-2 Survival Curve of an Infrastructure Facility**

### **2.1.3 Discrete Choice Model**

Another probabilistic model, discrete choice model, has also been used in the infrastructure performance modeling (see Zhang and Gao (2010), for example). Discrete choice models try to analyze choices between two or more discrete alternatives, such as staying or not staying in a certain condition state. Discrete choice model is similar to reliability model, but different in the number of condition states. Instead of defining only two condition states, discrete choice model allows the existence of multiple condition states.

One of the discrete choice model applications in infrastructure performance modeling can be explained as follows. Let  $C_n$  as the dependent variable represent the condition state for facility  $n$  and an underlying response variable  $U_n$  be a measure of the latent deterioration propensity for facility  $n$ .  $U_n$  is assumed as a continuous variable varying from  $-\infty$  to  $+\infty$ . The observed facility condition state  $k$  is a reflection of the latent variable  $U_n$ , which is specified to be a summation of a deterministic function of explanatory variables. In this case, the structure of the model can be described as:

$$U_n = \beta' X_n + \varepsilon_n \quad (n = 1, 2, \dots, N) \quad (2.6)$$

where  $U_n$  is the underlying response variable;  $X_n$  is a set of explanatory variables;  $\beta$  is the estimated parameter; and  $\varepsilon_n$  is the error term. The above equation cannot be directly estimated, since  $U_n$  is not observable. But the observable state  $k$  that facility  $n$  falls in can be used to estimate the parameters in the model. As such,  $C_n$  is governed by  $\Psi_k$ , the threshold values of the underlying response variable  $U_n$ . If the latent variable falls between the threshold  $\Phi_k$  and  $\Phi_{k-1}$ , then the  $C_n$  falls into the corresponding state  $k$ . In this regard, the thresholds separate the continuous underlying response variable  $U_n$  into different states. If  $\varepsilon_n$  is assumed to follow a standard normal distribution with mean 0 and standard deviation 1. Then the probability for facility  $n$  to be in the condition state  $k$  can be obtained by

$$P(C_{nk} = 1) = \Phi(\Psi_k - \beta' X_n) - \Phi(\Psi_{k-1} - \beta' X_n) \quad (2.7)$$



#### 2.1.4 Deterministic Models

Deterministic models are usually used for single facility performance modeling. Deterministic performance models can usually be expressed in a general form (2.8). Let  $\mathbf{y} = (y_1, \dots, y_n)$  be a  $n \times 1$  vector of sample of  $n$  condition observations expressed as:

$$y_i = h(\mathbf{x}_i, \boldsymbol{\beta}) + \varepsilon_i \quad (2.8)$$

where  $h$  is the deterioration function.  $\mathbf{x}_i$  is a  $1 \times p$  vector of  $p$  explanatory variables and  $\boldsymbol{\beta}$  a  $p \times 1$  vector of the corresponding coefficients. The error term  $\varepsilon_i$  is assumed to follow a certain distribution with associated coefficients  $\boldsymbol{\theta}$ . A data point affected by maintenance intervention can be modeled as:

$$y_i = h(\mathbf{x}_i, \boldsymbol{\beta}) + \delta_i A_i + \varepsilon_i \quad (2.9)$$

where  $\delta_i$  is an independent Bernoulli trial with success probability  $\tau$  representing the existence of the maintenance intervention.  $A_i$  represents maintenance effectiveness and is assumed to follow a probability distribution  $g$  with parameters  $\boldsymbol{\kappa}$ . Based on (2.9), the deterioration rate is captured by estimating the parameter  $\boldsymbol{\beta}$ . Moreover, by estimating the parameters  $\delta_i$  and  $A_i$ , the model determines if the  $i^{\text{th}}$  observation is affected by maintenance treatment

and if so, the magnitude of the impact. Details of the model can be found in Gao et al. (2011) and Hong and Prozzi (2010).

## **2.2 MAINTENANCE SCHEDULING MODELS**

Maintenance scheduling models try to find the optimal balance between costs and benefits of maintenance treatments and the most appropriate time to execute maintenance. Parameters considered in the scheduling include the cost of failure, the cost of preventive maintenance and the cost of rehabilitation and reconstruction. The foundation of any maintenance scheduling model relies on the underlying deterioration process and failure behavior of the facility. Maintenance scheduling optimization is one of the most critical issues in infrastructure asset management since the failure of a system during actual operation can be a costly and dangerous event. When a facility fails to operate in a system, it does not only cause damage to the system but also affect all the users.

Numerous efforts have been made to develop mathematical models as maintenance strategy decision-making aids. The infrastructure maintenance scheduling problem can generally be formulated in both discrete-time and continuous-time settings. In discrete-time setting models, a set of time points at which the maintenance treatment might be applied is predefined, for example, at the beginning or end of each year. The solution of this type of model determines which maintenance treatment should be applied at specific time points. In practice, infrastructure agencies make maintenance decisions subject to

budgetary constraints and resource availability. Agencies allocate resources for maintenance activities at the beginning of each budgeting year. It is therefore realistic to discretize the planning horizon into predetermined temporal stages (e.g., years) and restrict treatments to occur only at such time points. In continuous-time setting models, however, there is no predefined constraint about the timing of the maintenance treatment. The solution of this type of model determines both timing and type of maintenance treatments.

In the rest of this section, some popular maintenance scheduling models are discussed.

### **2.2.1 Markov Chain Based Linear Programming (LP)**

The Markov Chain based linear programming model is a discrete-time setting model. It is usually used for network-level infrastructure maintenance scheduling problem. In the LP model, facilities with similar deterioration patterns are grouped together. The solution of this model determines the percentage of a group's maintenance strategy instead of the strategy for each facility. Therefore, the computational effort of this type of model is simpler than the Integer Programming models (see 2.2.2) (Smilowitz and Madanat 2000; Guignier and Madanat 1999; Robelin and Madanat 2006; Wu et al. 2009; Gao et al. 2010).

The mathematical expression of the LP model can be explained as follows. Consider an infrastructure system as a set  $S = \{1, 2, \dots, S\}$  of different groups of facilities with homogeneous properties, e.g., by highway functional class.  $J = \{1, 2, \dots, J\}$  is defined as a set of state space with elements representing the

facility condition state. Each element of this set represents a specific condition state. In each time period, a decision should be made to determine the proportion of system that should receive maintenance treatment and the type of treatment that should be applied. A set of basic maintenance treatments is defined as  $\mathcal{M} = \{1, 2, \dots, M\}$ , where the  $M$ th treatment is set to be the most effective and also expensive. The scheduling time horizon is represented by the discrete set of time periods  $\mathcal{T} = \{1, 2, \dots, T\}$ .

Linear programming can be efficiently solved by the simplex algorithm. The simplex algorithm solves Linear Programming problems by constructing a feasible solution at a vertex of the polytope and then walking along a path on the edges of the polytope to vertices with non-decreasing values of the objective function until an optimum is reached. In general, the simplex algorithm is very efficient and can be guaranteed to find the global optimum if certain precautions against cycling are taken. Multiple optimal solutions are also possible in Linear programming problems. More details of the Markov Chain based linear programming model will be discussed in Chapter 3.

### **2.2.2 Integer Programming (IP)**

The IP model is another discrete-time setting approach to solving multiple facilities maintenance scheduling problem with budget constraints. The advantage of IP over LP is that its solution will assign maintenance treatments directly to individual facilities. However, it is usually applied to small-scale systems because the computational burden of combinatorics. Wang et al. (2003)

developed a multi-objective IP model for network-level pavement maintenance scheduling. The authors used the branch and bound algorithm to solve the proposed model. Ouyang and Madanat (2004) also developed an IP model outlining the scheduling of rehabilitation activities for multiple pavement facilities. They proposed a greedy heuristic method to solve the problem. However, due to the combinatory nature of the IP approach, the computational burden of network-level maintenance scheduling problem increases exponentially as the number of facilities under consideration increases. Therefore, some researchers tend to use approximation techniques when dealing with large-scale infrastructure maintenance scheduling. For example, Gao and Zhang (2008) use the approximate dynamic programming method to solve network-level pavement scheduling problem. Karabakal et al. (1994) and Dahl et al. (2008) use Lagrangian relaxation techniques to decompose the network-level IP problem into simpler project-level dynamic programming problems.

The mathematical formulation of the IP approach can be explained as follows. Let  $\mathcal{T} = \{1, 2, \dots, T\}$  represent the set of planning horizon.  $\mathcal{A}$  is defined as a set with  $N$  elements representing facilities in the system. A set of basic maintenance treatments is defined as  $\mathcal{M} = \{1, 2, \dots, M\}$ , where the  $M$ th treatment is set to be the most effective and expensive. Given the initial condition of facility  $a$ ,  $s_a^0$ , and the deterioration function  $f(\cdot)$ , the IP formulation is:

$$\max \sum_{t \in \mathcal{T}} \sum_{a \in \mathcal{A}} s_a^t \quad (2.10)$$

$$\text{s. t. } \sum_{a \in \mathcal{A}} \sum_{m \in \mathcal{M}} c_{amt} u_{amt} \leq B_t, \forall t \in \mathcal{T} \quad (2.11)$$

$$s_a^t = f(s_a^{t-1}) + \sum_{m \in \mathcal{M}} u_{amt} e_m, \quad \forall a \in \mathcal{A}, t \in \mathcal{T} \quad (2.12)$$

$$\sum_{m \in \mathcal{M}} u_{atm} \leq 1, \quad \forall a \in \mathcal{A}, t \in \mathcal{T} \quad (2.13)$$

$$u_{atm} \in \{0,1\}, \forall a \in \mathcal{A}, t \in \mathcal{T}, m \in \mathcal{M} \quad (2.14)$$

$$s_a^t > 0, \forall a \in \mathcal{A}, t \in \mathcal{T} \quad (2.15)$$

where,

$c_{amt}$  = maintenance cost of applying the  $m$ th treatment to facility  $a$  at year  $t$ ;

$B_t$  = budget at year  $t$ ;

$u_{amt}$  = binary variable, equals to 1 if the  $m$ th treatment is applied to facility  $a$  and equals to 0 otherwise;

$s_a^t$  = condition of facility  $a$  at year  $t$ ;

$f(\cdot)$  = deterioration function;

$e_m$  = maintenance effectiveness of the  $m$ th treatment.

The objective function (2.10) is to maximize the average annual condition of all facilities. Constraint (2.11) states that the annual expenditure cannot exceed the available budget. Constraint (2.12) represents the deterioration process of the infrastructure facility. Constraint (2.13) states that only one

treatment can be applied to the same facility each year. Constraints (2.14) and (2.15) define the decision variables of the IP model.

Methods of solving MIP problems can be largely classified into the following categories:

1. Branch and Bound method (BB). BB is the most widely used method for solving MIP problems. Subproblems are created by adding constraints to the integer variables. For example, if the  $k$ th integer variable  $x_k$ 's current solution is  $u$ , which is not an integer. Then the original problem is divided into two subproblems with respect to  $x_k$ , with  $x_k \leq \lfloor u \rfloor$  and  $x_k \geq \lceil u \rceil$  respectively. Lower bounds are obtained by the linear-programming relaxation to the problem. It is implemented by keeping the objective function and all constraints, but relaxing the integrality constraints. If the optimal solution to a relaxed problem is integral, it is an optimal solution to the subproblem, and the value can be used to terminate searches of subproblems whose lower bound is higher.
2. Branch and Cut method (BC). Branch-and-cut methods are exact algorithms for integer programming problems. It solves the integer programming problem by using a combination of cutting plane method with a branch-and-bound algorithm. It works by solving a sequence of linear programming relaxations of the integer programming problem. Cutting plane methods improve the linear relaxation of the problem to approximate integer solutions. The branch-and-bound algorithms proceed

by a sophisticated divide and conquer approach to solve problems. A Branch-and-cut algorithm can be described as follows.

Step 1. Initialization: Denote the original integer programming problem as the root node and store it in the waiting node list  $List$ . Set the upper bound to be  $UB := +\infty$  {best found}, the lower bound to be  $LB := -\infty$  {Best Possible} and current best solution  $\mathbf{x}^* := \emptyset$ . Go to step 2.

Step 2. Termination: If  $List = \emptyset$ , then the current best solution  $x^*$  which yielded the objective value  $obj$  is optimal; if  $\mathbf{x}^* = \emptyset$ , then the original problem is infeasible. If  $List \neq \emptyset$ , go to step 3.

Step 3. Node selection: Select and delete a node from  $List$ . Go to step 4.

Step 4. Relaxation: Solve the linear programming relaxation of the selected node problem. If the relaxation is infeasible, node is deleted. If an optimal integer solution  $\mathbf{x}_r$  is found and  $obj < UB$ , set  $UB := obj$ ,  $\mathbf{x}^* := \mathbf{x}_r$ , remove nodes  $j$  from  $List$  with  $LB_j > UB$  and go to step 2. If an optimal integer solution  $\mathbf{x}_r$  is found and  $obj \geq UB$ , remove nodes  $j$  from  $List$  with  $LB_j > UB$  and go to step 2. If the optimal solution  $\mathbf{x}_r$  is not integer, go to step 5.

Step 5. Add cutting planes: search for cutting planes that are violated by  $\mathbf{x}_r$ ; if any are found, add them to the relaxation and go to step 4. If no cutting planes are found, go to step 6.



Step 6.        Branching: find variable  $x_k$  of solution  $\mathbf{x}_r$  with fractional value  $v$ . Create node  $j_{new}$  with bound  $x_k \leq [v]$  and set  $LB_{j_{new}} := obj$ . Store node  $j_{new}$  in  $List$ . Create node  $j_{new}$  with bound  $x_k \geq [v]$  and set  $LB_{j_{new}} := obj$ . Store node  $j_{new}$  in  $List$ . Go to step 2.

3. Branch and Price method (BP). This is essentially branch and bound combined with column generation. This method is used to solve integer programs where there are too many variables to represent the problem explicitly. Thus only the active set of variables is maintained and columns are generated as needed during the solution of the linear program.

### 2.2.3 Reliability Model

Reliability model is one of the continuous-time setting models. It is usually used to model maintenance plan for infrastructure facilities whose failures can be clearly defined (e.g., bridge, traffic lights).

For a single facility, there are different reliability models that can be used, including age replacement models, minimal repair models, and inspection/maintenance models. Age replacement models deal with optimal replacement policies, which are based on age dependent operating costs. Minimal repair models focus on repairing a failed unit rather than replacing it. They usually combine a periodic replacement policy with a minimal repair activity upon a unit failure. Finally, inspection/maintenance models are concerned with

maintenance policies in which the current state of a system is not known but is available through an inspection. For multi-facility systems, reliability maintenance models aim at optimal maintenance policies for a system consisting of several units of facilities, which may or may not depend on each other. Multi-facility reliability maintenance models include block or group maintenance models, inventory models and opportunistic models.

One example of the reliability model application is to develop an optimal replacement policy that will minimize the sum of operating and replacement costs per unit time. The replacement policy is to perform replacements at time intervals of length  $t_r$ . The objective is to determine the optimal interval between replacements to minimize the total cost of operation and replacement per unit time. The total cost per unit time, for replacement at time  $t_r$ , can be expressed as

$$C(t_r) = c(t) + C_r = \frac{1}{t_r} \left[ \int_0^{t_r} c(t) dt + C_r \right] \quad (2.16)$$

where,

$c(t)$  = operating cost per unit time at time  $t$  after replacement;

$C_r$  = cost of a replacement.

#### **2.2.4 Optimal Control Model**

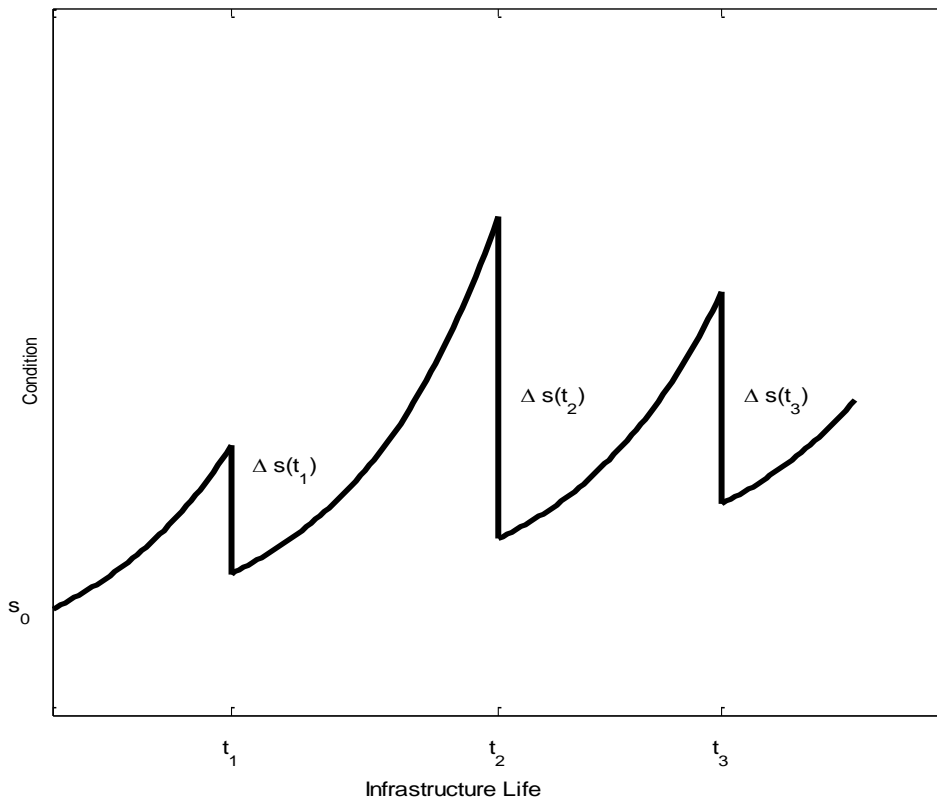
The optimal control model is a continuous-time setting approach to modeling project-level maintenance scheduling problems. The optimal control model deals

with the problem of finding a control law for a given system such that a certain optimality criterion is achieved. A control problem includes a cost functional that is a function of state and control variables. An optimal control is a set of differential equations describing the paths of the control variables that minimize the cost function. The optimal control solution can be derived using Pontryagin's maximum principle or by solving the Hamilton-Jacobi-Bellman equation.

The advantage of optimal control model is that maintenance actions are not restricted at fixed time points. The solution of the model determines the optimal maintenance timing. However, unlike ordinary optimal control problems, the control actions (maintenance treatments) in infrastructure maintenance scheduling problem are impulsive, leading to sudden jumps in the facility's condition trajectory. Therefore, special techniques are required to solve this type of problem. Researchers have adopted different approaches to addressing this issue in previous studies. For example, Tsunokawa and Schofer (1994) developed an approximation method for road maintenance scheduling problems. In their paper, the impulse control problem is simplified by approximating discrete controls using ordinary continuous controls. The simplified problem can then be solved effectively using the Pontryagin's maximum principle. In another paper, Li and Madanat (2002) solved the same problem using the assumption that the planning horizon is infinite and the condition of the facility will enter a steady state after the first maintenance treatment. Under this assumption, the optimal resurfacing strategy is based on a minimum serviceability level. When the facility condition deteriorates to that level, the condition should be brought back to the best serviceability level. Ouyang and Madanat (2006) derived an exact analytical

solution for the same resurfacing problem with a finite planning horizon. Through variational derivation, the author obtained the necessary condition of the control problem.

As described by Tsunokawa and Schofer (1994), the condition of an infrastructure facility (e.g., International Roughness Index), denoted by  $s$ , usually follows a special trajectory curve over time as the facility deteriorates and receives maintenance treatments as shown in Figure 2-3.



**Figure 2-3 Infrastructure Condition Trend Trajectory**

The deterioration rate of the facility is assumed to be a function,  $f$ , of the current condition level and expressed as follows:

$$\dot{s}(t) = f(s(t)) \quad (2.17)$$

The amount of condition improvement after a maintenance treatment is assumed to be a function,  $g$ , of the maintenance intensity (e.g., thickness of overlay),  $w$ , and the condition level immediately before the treatment, written as:

$$\Delta s(t) = g(s(t), w) \quad (2.18)$$

The initial condition is expressed as:

$$s(0) = s_0 \quad (2.19)$$

Costs for the agency and user are assumed to be functions of the condition and the maintenance intensity, respectively, and are written as  $C(s)$  and  $M(w)$ .

Using these functions, the total life-cycle costs for the agency and user of an infrastructure facility can be written as follows:

$$J = \int_0^T C(s(t)) dt + \sum_{i=1}^N M(w(t_i)) \quad (2.20)$$

In this formula,  $T$  is the planning horizon,  $i$  represents the  $i^{th}$  maintenance action, and  $N$  is the number of maintenance actions during the horizon. The problem of finding the optimal maintenance strategy for a given infrastructure facility can be solved by evaluating the optimal values of  $t_i$  and  $w(t_i)$ , which minimize the life cycle cost. The condition trajectory,  $s(t)$ , is determined by Eqs. (2.17) and (2.18) and the initial condition (2.19). Therefore, the problem introduced above can be solved as an impulse control problem. This type of problem arises frequently in practice. There are two basic approaches to the solution. One relies on unconventional quasi-variational inequalities and uses the dynamic programming methodology. The other approach, similar to the classical calculus of variations, formulates the optimality necessary conditions in terms of maximum principle.

### **2.3 INFRASTRUCTURE MAINTENANCE SCHEDULING MODELS CONSIDERING BUDGET UNCERTAINTY**

The aforementioned models treat the annual budget as a fixed amount. An underlying assumption is that the actual funds to support maintenance activities never deviate from the original expectation. Under this assumption, the maintenance scheduling problem can be solved optimally using the optimization models discussed previously. However, this assumption is often unrealistic, because the funding allocated to infrastructure maintenance activities is usually subject to uncertainties due to various financial and political risks. Consequently, the actual amount of money distributed to maintenance activities may deviate

from the original estimate. If funding falls short for some of the years during the planning period, part of the planned maintenance activities may be suspended or postponed, leading to inevitable condition fluctuation for the whole system. Therefore, ignoring the random characteristics of the future budget may limit the usefulness of the optimal maintenance scheduling solution.

In recent years, several researchers addressed the problems of budget uncertainty in the infrastructure management area. For example, Li and Puyan (2006) formulated a highway project selection problem under budget uncertainty as a multi-choice multidimensional Knapsack problem with multi-stage budget recourses. In their paper, the objective is to select a subset of candidate projects to achieve maximized system benefits under budget and other constraints. Gao and Zhang (2008) investigated the uncertainties in the pavement deterioration process and proposed a robust optimization approach for project-level maintenance planning problem. Using this approach, the decision maker is able to control the probability of achieving a certain level of condition requirement by adjusting the amount of money invested. Wu and Flintsch (2009) proposed a chance-constrained programming model with the ability to control the probability of going over budget for network-level facility maintenance planning problems. The solution of the proposed model is obtained by first choosing a conservative value for the budget and then treating the budget as fixed. However, the obtained scheduling solution of this model is only optimal at a given probability.

In this research, the network-level infrastructure maintenance scheduling problem under budget uncertainty is formulated as a multi-stage, linear stochastic programming model. Stochastic programming is a framework for

modeling optimization problems that involve uncertainty. The goal of stochastic programming is to find a solution that is feasible for all data scenarios. Stochastic programming models take advantage of the fact that probability distributions governing the data are known or can be estimated. In this research, the proposed stochastic programming approach differs from its deterministic counterpart in that it attempt to achieve the best expected objective value over all possible realizations of the random budgets.



## CHAPTER 3      PROBLEM FORMULATION

The mathematical formulation of the model developed in this research is presented in this section. The performance model is first discussed. Then the deterministic version of the infrastructure maintenance scheduling problem is presented. Finally, the author introduces the stochastic extension of the deterministic formulation by taking budget uncertainty into account.

### 3.1 NOTATIONS

The sets, parameters, and variables used in the model and algorithm development are presented as follows.

**Table 3.1 Notation**

<b>Sets</b>	
<b>S</b>	set of facility groups and $\mathbf{S} = \{1, 2, \dots, S\}$
<b>I</b>	set of facility condition states and $\mathbf{I} = \{1, 2, \dots, I\}$ with $I$ represents the worst condition state
<b>M</b>	set of maintenance treatments and $\mathbf{M} = \{1, 2, \dots, M\}$ with the $M^{th}$ treatment being the most effective and expensive
<b>T</b>	set of planning periods $\mathbf{T} = \{1, 2, \dots, T\}$
<b>K</b>	set of all nodes in the scenario tree and $\mathbf{K} = \{1, \dots, K\}$ , where $k = 1$

	corresponds to the root node at $t=1$ and $t(k)$ denotes the year corresponding to node $k$
$\mathbf{N}$	set of scenarios and $\mathbf{N} = \{1, 2, \dots, N\}$
<b>Parameters</b>	
$B_t$	available budget at time period $t$
$\tilde{B}_t$	random variable representing available budget at time period $t$
$B_t^n$	realization of the budget random variable $\tilde{B}_t$ in scenario $n$
$b_t$	number of realizations of $\tilde{B}_t$
$C_{smt}$	unit cost of applying the $m^{th}$ treatment to the $s^{th}$ facility group at time period $t$
$L_s$	number of the $s^{th}$ facility group
$P_{sijm}$	deterioration transition probability from condition state $i$ to state $j$ when the $m^{th}$ treatment is applied to the $s^{th}$ facility group. $P_{sijm}$ satisfies the constraint of $\sum_{j \in \mathbf{I}} P_{sijm} = 1, s \in \mathbf{S}, (i, j) \in \mathbf{I}, m \in \mathbf{M}$
$X_{si1}$	proportion of the $s^{th}$ facility group in condition state $i$ at the beginning of the first time period, which is known to the decision maker before the maintenance planning
$X^*$	minimum requirement on the proportion of facilities in the first condition state
$P(t)$	condition state probability vector of a facility at time period $t$
$p_i(t)$	probability that a facility stays in state $i$ in time period $t$ ; $i \in \mathbf{I}$ , larger value of $i$ corresponds to worse condition state; $\sum_{i=1}^n p_i(t) = 1$

$D$	transition probability matrix
$d_{ij}$	probability that the facility will deteriorate from state $i$ to state $j$ in one time period, if $i < j$ and $(i, j) \in \mathbf{I}$ ; probability that the facility will stay in the same state in one time period, if $i = j$ and $(i, j) \in \mathbf{I}$
$p^n$	probability of occurrence of the $n^{th}$ scenario and $\sum_{n \in \mathbf{N}} p^n = 1$
<b>Variables</b>	
$X_{simt}$	proportion of the $s^{th}$ facility group in condition state $i$ that receives the $m^{th}$ treatment at time period $t$
$X_{simt}^n$	decision variable $X_{simt}$ for scenario $n$
$X_{sit}$	proportion of the $s^{th}$ type facility in condition state $i$ at time period $t$
$M_{smt}$	proportion of the $s^{th}$ facility group that receives the $m^{th}$ treatment at time period $t$

### 3.2 DETERIORATION MODELING

The concept of infrastructure condition is developed to quantitatively relate the condition of a facility to its ability to serve its users. Infrastructure condition is often represented by discrete ratings or states. Using discrete ratings instead of continuous indicators simplifies the computational complexity of the maintenance decision-making process, as details are not necessary at this level of management. In this research, the deterioration of a facility is modeled as a discrete-time, state-based model widely used in infrastructure management (Wang et al. 1994; Li et al. 1996; Abaza et al. 2004; Gao and Zhang 2007).

The basic idea of the discrete-time, state-based model is introduced as follows. Facility condition at different years is represented by a condition state probability vector:

$$P(t) = [p_1(t), \dots, p_I(t)]^T \quad (3.1)$$

The deterioration process of a facility can be expressed by the change of the elements in the condition state probability vectors  $P(t)$ . A transition probability matrix  $D$  can be used to calculate this change.

$$D = \begin{pmatrix} d_{11} & d_{12} & \dots & d_{1I} \\ 0 & d_{22} & \dots & d_{2I} \\ \vdots & \vdots & & \vdots \\ 0 & 0 & \dots & 1 \end{pmatrix} \quad (3.2)$$

Because a facility cannot improve to a better condition state by itself, the elements  $d_{ij}$  is replaced by 0 for  $i > j$ . Furthermore, the value of 1 in the last row of  $D$  corresponding to state  $I$  indicates that the condition cannot deteriorate further. From all the above, the future condition can be predicted as:

$$P(t+1) = D \times P(t) \quad (3.3)$$

where,  $P(t+1)$  represents the condition state probability vector at time  $t+1$ ;

### 3.3 MAINTENANCE PLANNING MODEL WITH DETERMINISTIC BUDGETS

Consider an infrastructure system as a set  $\mathbf{S}$  of facilities, e.g., pavements, bridges, rail, mass transit, and dams. Condition  $\mathbf{I} = \{1, 2, \dots, I\}$  is defined as a set of state space with elements representing the facility condition in which 1 represents the best condition state and  $I$  the worst. A set of basic maintenance treatments is defined as  $\mathbf{M} = \{1, 2, \dots, M\}$ , where the  $M^{th}$  maintenance treatment is set to be most effective and expensive. The scheduling time horizon is represented by the discrete set of time periods  $\mathbf{T} = \{1, 2, \dots, T\}$ . During each time period, the conditions of the facilities deteriorate due to usage, aging, and environment. The maintenance treatment applied at time period  $t$  will affect the condition at time period  $t+1$ .

Using the discrete-time, state-based deterioration model, the infrastructure maintenance scheduling problem with deterministic budgets is formulated in Equations (3.4)-(3.8).

$$\max \frac{1}{\sum_{s \in \mathbf{S}} L_s} \frac{1}{T+1} \left( \sum_{t=1}^T \sum_{s=1}^S \sum_{m=1}^M L_s X_{s1mt} + \sum_{s=1}^S \sum_{i=1}^I \sum_{m=1}^M P_{i1m} L_s X_{simT} \right) \quad (3.4)$$

$$\text{s.t. } \sum_{m=1}^M X_{sim1} = X_{si1}, \forall s \in \mathbf{S}, i \in \mathbf{I} \quad (3.5)$$

$$\sum_{m=1}^M X_{sjmt} = \sum_{m=1}^M \sum_{i=1}^I P_{sijm} X_{sim,t-1}, \forall s \in \mathbf{S}, j \in \mathbf{I}, t = 2, \dots, T \quad (3.6)$$

$$\sum_{s=1}^S \sum_{i=1}^I \sum_{m=1}^M C_{smt} X_{simt} L_s \leq B_t, \forall t \in \mathbf{T} \quad (3.7)$$

$$0 \leq X_{simt} \leq 1, \forall s \in \mathbf{S}, i \in \mathbf{I}, m \in \mathbf{M}, t \in \mathbf{T} \quad (3.8)$$

The objective (3.4) of the planning problem is to maximize the proportion of all facilities in the best condition state over the planning horizon. The first term inside the parenthesis represents the proportion from time period 1 to time period  $T$ . The second term in the parenthesis represents the proportion at time period  $T + 1$ , because a facility's condition at time period  $T + 1$  is fully determined by its condition and applied maintenance treatments at time period  $T$ . Constraint (3.5) represents the initial condition of each facility group at the beginning of the planning horizon. Constraint (3.6) represents the deterioration process of the facilities between two consecutive time periods. Constraint (3.7) ensures that the annual expenditure of maintenance activities does not exceed the budget. Once the decision variables  $X_{simit}$  of problem (3.4)–(3.8) are obtained, the condition of each facility group can be calculated as:

$$X_{sit} = \sum_{m=1}^M X_{simit}, \forall s \in \mathbf{S}, i \in \mathbf{I}, t \in \mathbf{T} \quad (3.9)$$

The maintenance decision is then calculated as:

$$M_{smt} = \sum_{i=1}^I X_{simit}, \forall s \in \mathbf{S}, m \in \mathbf{M}, t \in \mathbf{T} \quad (3.10)$$

### 3.4 MAINTENANCE PLANNING MODEL UNDER BUDGET UNCERTAINTY

In this research, the budget uncertainty in the infrastructure maintenance scheduling problem is modeled using the Stochastic Programming model. Stochastic Programming is a framework for modeling optimization problems that involve uncertainty. Stochastic Programming model takes advantage of the fact that probability distributions governing the data are known or can be estimated. The objective of using Stochastic Programming is to find the solution that is feasible for all the possible data scenarios and maximize (or minimize) the expectation of some functions of the decisions and the random variables.

In a multi-period infrastructure maintenance scheduling problem, the budgets in time period 2 to  $T$  are unknown to decision makers at period 1. Therefore, to extend the deterministic formulation (3.4)-(3.8) to a stochastic setting, the budget  $B_t$  at time period  $t$ ,  $t=1, \dots, T$  is replaced with a random variable  $\tilde{B}_t$ . In this research,  $\tilde{B}_t$  is assumed to evolve as a discrete time stochastic process with a finite probability space represented in the form of a scenario tree (for example Figure 3-1).

The  $T$  stages in the tree represent  $T$  planning periods. The nodes at stage  $t$  of the tree correspond to scenarios of possible values of  $\tilde{B}_t$ . If  $b_t$  represents the number of realizations of  $\tilde{B}_t$ , then there are  $\prod_{t=1}^T b_t$  nodes at the  $T$ th stage of the tree.

Furthermore, let  $\mathbf{K} := \{1, \dots, K\}$  denote the set of all nodes, where  $k=1$  corresponds to the root node at  $t=1$  and  $t(k)$  denotes the time period corresponding to node  $k$ . Each node  $k$  is connected to its parent node  $k^+$  at

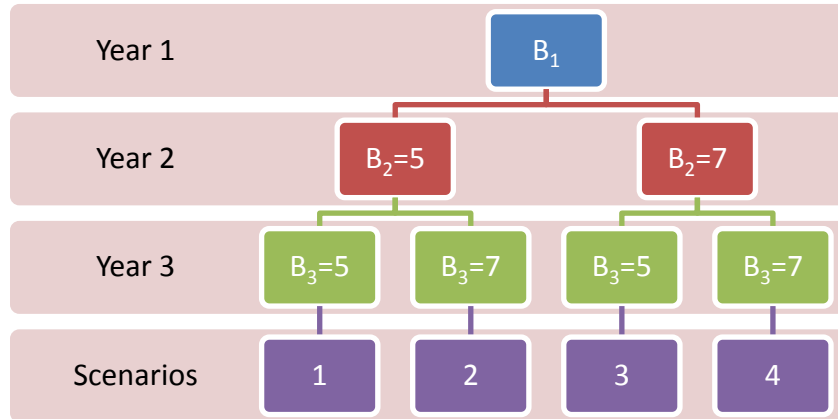
time period  $t-1$  by an arc. A set of child nodes is associated with each node  $k$  with  $t(k) \in \{1, \dots, T-1\}$ . The node set  $(1, \dots, (k^+)^+, k^+, k)$  is defined as a path from the root to node  $k$ .

A set  $\mathbf{N} = \{1, 2, \dots, N\}$  is defined as the scenarios with each element representing a path from the root to any nodes  $k$  with  $t(k) = T$ . A scenario represents one possible combination of values for all uncertain budgets. The probability associated with a scenario is the probability of reaching the corresponding node at year  $T$  from the root node. For each scenario, the associated probability is  $p^n$  and  $\sum_{n \in \mathbf{N}} p^n = 1$ .

To illustrate the concept of the scenario tree, a simple example is presented with a planning period of three years (Figure 3.1). The budget at the starting time period is already known to the decision maker. It is assumed that there are two possible values, \$5 million and \$7 million, for both the second and third year budgets. Therefore, four possible scenarios ( $N = 4$ ) may occur over the three decision periods. With the scenarios defined above, a probability of  $p^n = 0.25, n \in \{1, 2, 3, 4\}$  is assigned to each scenario. For a real problem, the decision maker can assign any probability to each scenario based on his/her own judgment.

As illustrated in Figure 3-1, the scenario tree is divided into branches corresponding to different realizations of the budget random variable. For example, the budget at year 2 is \$5 million for scenarios 1 and 2 and \$7 million for scenarios 3 and 4. For scenarios 1 and 3, the budget at year 3 is \$5 million, while for scenarios 2 and 4, the budget at year 3 is \$7 million.





**Figure 3-1 Scenario Tree**

If scenarios  $n_1, n_2$  ( $n_1, n_2 \in \mathbf{N}$ ) have the same information state at time period  $t$  (sharing the same node at  $t$  in the scenario tree), the two scenarios are indistinguishable at  $t$ . In general, scenarios  $n_1, n_2$  are indistinguishable at  $t$  if they are identical in realizations for all uncertain budgets up to time  $t$ . For example, in Figure 3-1, scenarios 1 and 2 are indistinguishable at year 2, as they have the same budget realization at that year. However, they are distinguishable at year 3, because their budgets at that year are different. Moreover, let  $\mathbf{t}(n_1, n_2)$  denote the latest time period at the end of which scenarios  $n_1$  and  $n_2$  are indistinguishable. For example, in Figure 3-1,  $\mathbf{t}(1, 2) = 2$ , scenarios 1 and 2 differ in terms of budget realization after year 2.

Using the notations discussed above, the infrastructure maintenance scheduling problem under budget uncertainty can be formulated as a multi-stage stochastic programming problem with Equations (3.11)-(3.16).

$$\max \frac{1}{\sum_{s \in \mathbf{S}} L_s} \frac{1}{T+1} \left( \sum_{n=1}^N p^n \left( \sum_{t=1}^T \sum_{s=1}^S \sum_{m=1}^M L_s X_{s1mt}^n + \sum_{s=1}^S \sum_{i=1}^I \sum_{m=1}^M P_{si1m} L_s X_{simT}^n \right) \right) \quad (3.11)$$

$$\text{s.t. } \sum_{m=1}^M X_{sim1}^n = X_{si1}, \forall s \in \mathbf{S}, i \in \mathbf{I}, n \in \mathbf{N} \quad (3.12)$$

$$\sum_{m=1}^M X_{sjmt}^n = \sum_{m=1}^M \sum_{i=1}^I P_{sijm} X_{sim,t-1}^n, \forall s \in \mathbf{S}, j \in \mathbf{I}, t = 2, \dots, T, n \in \mathbf{N} \quad (3.13)$$

$$\sum_{n=1}^N \sum_{s=1}^S \sum_{i=1}^I \sum_{m=1}^M C_{smt} X_{simt}^n L_s \leq B_t^n, \forall s \in \mathbf{S}, j \in \mathbf{I}, t \in \mathbf{T}, n \in \mathbf{N} \quad (3.14)$$

$$0 \leq X_{simt}^n \leq 1, \forall s \in \mathbf{S}, i \in \mathbf{I}, m \in \mathbf{M}, t \in \mathbf{T}, n \in \mathbf{N} \quad (3.15)$$

$$X_{simt}^{n_1} = X_{simt}^{n_2}, \forall s \in \mathbf{S}, i \in \mathbf{I}, m \in \mathbf{M}, t \in \mathbf{T}, n_1 \in \mathbf{N}, n_2 \in \mathbf{N} \quad (3.16)$$

The objective function (3.11) maximizes the expected annual proportion of facilities in the best condition state over the probability space of random variable  $\tilde{B}_t$ . Constraints (3.12)–(3.15) are the same as constraints (3.5)–(3.8) but for the  $n^{\text{th}}$  scenario. Decisions for different scenarios are linked by nonanticipativity constraints (3.16). The nonanticipativity-constraint states that decision variables of scenario  $n_1$  and  $n_2$  are equal whenever  $n_1$  and  $n_2$  are indistinguishable. In Stochastic Programming, constraints enforcing such conditions are called nonanticipativity constraints, implying that the future cannot be anticipated. The nonanticipativity-constraint acts as a coupling constraint that connects different

scenarios together and specifies how the information on budget is shared among scenarios. For example, in Figure 3-1, decision variables before year 3 should be the same for scenarios 1 and 2.

## CHAPTER 4 SOLUTION PROCEDURE

Multi-stage stochastic programming is one of the most difficult problems in mathematical programming. The basic approach to multistage stochastic programs is to approximate the stochastic process using a process of finite scenarios exhibiting a tree structure. The size of the problem grows quickly as the number of stages and number of scenarios increase, typically leading to very large-scale linear programming models.

Existing computational methods for multistage stochastic programming problems include decomposition methods that exploit specific structures of the model to split it into manageable pieces and scenario reduction techniques that generate smaller scenario trees from an initial set of scenarios. Decomposition methods can be further classified into two groups: 1) primal decomposition methods that define subproblems according to time stages; and 2) dual methods that construct subproblems that correspond to scenarios.

In this research, the author proposes the use of the augmented Lagrangian decomposition method (Rosa and Ruszczyński 1996) and scenario reduction method (Heitsch and Romisch 2009). The major computational advantage of the augmented Lagrangian decomposition method is the possibility of solving the dual problem by the multiplier method. Another important advantage of the augmented Lagrangian decomposition method over the usual Lagrangian duality is its sufficiency for primal recovery when the dual solution is known. The advantage of the scenario reduction method is that it significantly

simplifies the computational effort. The following sections introduce the basic principle of these two methods.

#### 4.1 AUGMENTED LAGRANGIAN DECOMPOSITION (ALD)

Let  $X_1, X_2, \dots, X_L$  be non-empty closed convex subsets, and let  $f_i, i = 1, 2, \dots, L$ , be convex functions. Moreover, let  $A_i$  be matrices of dimension  $m \times n_i, i = 1, 2, \dots, L$ , and let  $b \in R^m$ . Consider the following convex programming problem:

$$\min \sum_{i=1}^L f_i(x_i) \quad (4.1)$$

$$\sum_{i=1}^L A_i x_i = b \quad (4.2)$$

$$x_i \in X_i, i = 1, 2, \dots, L \quad (4.3)$$

Problems (4.1)-(4.3) can be decomposed into  $L$  smaller and simpler problems  $\{\min_{x_i \in X_i} f_i(x_i), i = 1, \dots, L\}$  if constraint (4.2) is relaxed. To use this special structure to solve the problem, the augmented Lagrangian function is defined for this problem as:

$$\Lambda(\mathbf{x}, \pi) = \sum_{i=1}^L f_i(x_i) + \pi \left( b - \sum_{i=1}^L A_i x_i \right) + \frac{\rho}{2} \left\| b - \sum_{i=1}^L A_i x_i \right\|^2 \quad (4.4)$$

where  $\rho$  is the penalty parameter and  $\rho > 0$ . The dual problem is also defined as:

$$\max_{\pi \in R^m} \{g(\pi) = \inf_{\mathbf{x} \in X} \Lambda(\mathbf{x}, \pi)\} \quad (4.5)$$

For every optimal solution  $\hat{\pi}$  of (4.5), a point  $\hat{\mathbf{x}}$  is a solution of (4.1)–(4.3) only if  $\Lambda(\hat{\mathbf{x}}, \hat{\pi}) = \min_{\mathbf{x} \in X} \Lambda(\mathbf{x}, \hat{\pi})$ . Therefore, the optimal solution of problems (4.1)–(4.3) is obtained by solving the dual problem (4.5) instead (Ruszczynski 1997). The dual problem is solved by iteratively using the method of multipliers (4.6)–(4.7) until a convergence is reached (Sun and Yuan 2006):

$$\mathbf{x}^k = \arg \min_{\mathbf{x} \in X} \Lambda(\mathbf{x}, \pi^k) \quad (4.6)$$

$$\pi^{k+1} = \pi^k + \rho(b - \mathbf{A}\mathbf{x}^k), k = 0, 1, 2, \dots \quad (4.7)$$

where  $k$  is the iteration counter for the method of multipliers.

Thus far, although the coupling constraint (4.2) is relaxed, solving (4.6) is still cumbersome, because the third term of (4.4) is inseparable. As a result, problem (4.6) cannot be split into smaller subproblems for  $x_i, i = 1, 2, \dots, L$ . To overcome this difficulty, an iterative nonlinear Jacobi method to the minimization of (4.4) is applied (Ruszczynski 1997; Rosa and Ruszczynski 1996). This method uses a certain approximation of the minimizer  $\mathbf{x}^k$  in (4.6) and solves the following simplified functions for  $i = 1, 2, \dots, L$ :

$$\Lambda_i(x_i, \tilde{\mathbf{x}}_i, \boldsymbol{\pi}) = f_i(x_i) - A_i^T \boldsymbol{\pi} x_i + \frac{\rho}{2} \left\| b - A_i x_i - \sum_{j \neq i} A_j x_j \right\|^2 \quad (4.8)$$

where  $\tilde{\mathbf{x}}_i$  represents all the solutions  $x_j$  with  $j \neq i$ . The main goal of this approach is to replace (4.6) with  $L$  smaller problems:

$$\min_{x_i \in X_i} \Lambda_i(x_i, \tilde{\mathbf{x}}_i, \boldsymbol{\pi}^k), i = 1, 2, \dots, L \quad (4.9)$$

and to iteratively update the parameter  $\tilde{\mathbf{x}}_i, i = 1, 2, \dots, L$ . In this sense, solving (4.9) is equivalent to solving (4.6) with respect to  $x_i$  while keeping all  $x_j, j \neq i$  fixed. In this way, (4.6) can be solved using the Jacobi method (Rosa and Ruszczynski 1996):

*Step 0:* Set the iteration counter of the Jacobi method  $r := 0$  and determine the initial solution values  $\mathbf{x}^{k,0} = \mathbf{x}^{k-1}$ .

*Step 1:* Set  $r := r + 1$ . Solve (5.9) for  $i = 1, 2, \dots, L$  and obtain the solution  $\mathbf{x}^{k,r}$ , where  $\mathbf{x}^{k,r} = \{x_1^{k,r}, x_2^{k,r}, \dots, x_L^{k,r}\}$ .

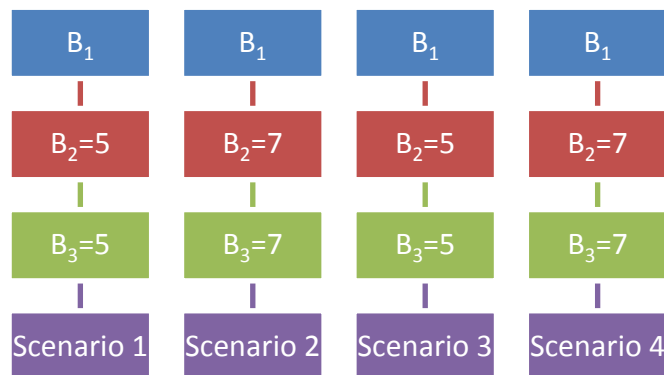
*Step 2:* If  $\mathbf{A}\mathbf{x}^{k,r} = \mathbf{A}\mathbf{x}^{k,r-1}$  then stop and set  $\mathbf{x}^k = \mathbf{x}^{k,r}$ ; otherwise update  $\mathbf{x}^{k,r}$  by (5.10) and go to Step 1:

$$\mathbf{x}^{k,r} = \mathbf{x}^{k,r-1} + \tau(\mathbf{x}^{k,r} - \mathbf{x}^{k,r-1}) \quad (4.10)$$

where  $\tau$  is a weighting factor.

## 4.2 APPLICATION OF ALD TO STOCHASTIC PROGRAMMING

Using the idea discussed in the previous section, an infrastructure maintenance scheduling problem under budget uncertainty (3.11)-(3.16) can be decomposed into  $N$  subproblems ( $N$  scenarios), if the nonanticipativity constraint (3.16) is relaxed. Using the Figure 4-1 as an example, if the nonanticipativity constraint is relaxed, the scenario tree will be separated to individual branches (Figure 4-1).



**Figure 4-1 Scenario Tree after Decomposition**

Because of the special structure of the problem, the augmented Lagrangian decomposition method can be used. The augmented Lagrangian function  $\Lambda$  is first defined as:



$$\begin{aligned}
\Lambda(x, \pi) = & \frac{1}{\sum_{s \in \mathbf{S}} L_s} \frac{1}{T+1} \left( \sum_{n=1}^N p^n \left( \sum_{t=1}^T \sum_{s=1}^S \sum_{m=1}^M L_s X_{s1mt}^n + \sum_{s=1}^S \sum_{i=1}^I \sum_{m=1}^M P_{i1m} L_s X_{simT}^n \right) \right) \\
& - \sum_{n=1}^N \sum_{n'=1}^N \sum_{t=1}^T \sum_{s=1}^S \sum_{i=1}^I \sum_{m=1}^M \pi_{stim}^{n,n'} (X_{stim}^n - X_{stim}^{n'}) + \frac{\rho}{2} \sum_{n=1}^N \sum_{n'=1}^N \sum_{t=1}^T \sum_{s=1}^S \sum_{i=1}^I \sum_{m=1}^M (X_{stim}^n - X_{stim}^{n'})^2
\end{aligned} \tag{4.11}$$

The subproblem of the  $n^{th}$  scenario is expressed as:

$$\begin{aligned}
\max & \frac{1}{\sum_{s \in \mathbf{S}} L_s} \frac{1}{T+1} p_s \left( \sum_{t=1}^T \sum_{s=1}^S \sum_{m=1}^M L_s X_{s1mt}^n + \sum_{s=1}^S \sum_{i=1}^I \sum_{m=1}^M P_{i1m} L_s X_{simT}^n \right) \\
& - \sum_{n=1}^N \sum_{n'=1}^N \sum_{t=1}^T \sum_{s=1}^S \sum_{i=1}^I \sum_{m=1}^M \pi_{stim}^{n,n'} X_{stim}^n + \frac{2}{\rho} \sum_{n=1}^N \sum_{n'=1}^N \sum_{t=1}^T \sum_{s=1}^S \sum_{i=1}^I \sum_{m=1}^M (X_{stim}^n)^2 \\
& - \frac{\rho}{2} \sum_{n=1}^N \sum_{n'=1}^N \sum_{t=1}^T \sum_{s=1}^S \sum_{i=1}^I \sum_{m=1}^M 2X_{stim}^n \tilde{X}_{stim}^{n'}
\end{aligned} \tag{4.12}$$

$$\text{s.t. } \sum_{m=1}^M X_{sim1}^n = X_{si1}, \forall s \in \mathbf{S}, i \in \mathbf{I}, \tag{4.13}$$

$$\sum_{m=1}^M X_{sjmt}^n = \sum_{m=1}^M \sum_{i=1}^I P_{sijm} X_{sim,t-1}^n, \forall s \in \mathbf{S}, j \in \mathbf{I}, t = 2, \dots, T \tag{4.14}$$

$$\sum_{n=1}^N \sum_{s=1}^S \sum_{i=1}^I \sum_{m=1}^M C_{smt} X_{simt}^n L_s \leq B_t^n, \forall s \in \mathbf{S}, i \in \mathbf{I}, t \in \mathbf{T} \tag{4.15}$$

$$0 \leq X_{simt}^n \leq 1, \forall s \in \mathbf{S}, i \in \mathbf{I}, m \in \mathbf{M}, t \in \mathbf{T} \tag{4.16}$$

The problem (4.12)-(4.16) is minimized with respect to decision variables associated with the  $n^{th}$  scenario assuming that decision variables of other

scenarios are temporarily fixed. As suggested by Rosa and Ruszczyński (1996), scenarios are numbered so that at the  $i^{th}$  scenario, the  $i + 1^{th}$  scenario has the largest last common stage with  $i$  among all scenarios  $j > i$ . The augmented Lagrangian decomposition algorithm is carried out in the order of 1 to  $N$ . By applying the method of multipliers and the Jacobi method, the infrastructure maintenance scheduling problem can be solved.

### 4.3 SCENARIO REDUCTION (SR)

Scenario reduction is about eliminating scenarios that are similar to other scenarios. For a given Stochastic Programming problem, a large number of scenarios usually exist. These scenarios normally result from a simulation where the distribution of the simulated random variable is known. The aim of the scenario reduction is that a reduced number of scenarios still represent the underlying distribution in an acceptable way.

Assume that the original probability distribution  $P$  is discrete and carried by a finite number of scenarios  $\omega_i \in \Omega$  with weights  $p_i > 0, i = 1, \dots, N$ , and  $\sum_{i=1}^N p_i = 1$ , i.e.,  $P = \sum_{i=1}^N p_i \delta_{\omega_i}$ . Let  $J \subset \{1, \dots, N\}$  and consider the probability measure  $Q$  having scenario  $\omega_j$  with probabilities  $q_j, j \in \{1, \dots, N\} \setminus J$ , i.e., compared to  $P$  the measure  $Q = \sum_{j \notin J} q_j \delta_{\omega_j}$  is reduced by deleting all scenarios  $\omega_j, j \in J$  and by assigning new probabilistic weights  $q_j$  to each scenario  $\omega_j, j \notin J$ .

One of the algorithms used for reducing scenarios is to delete one scenario at a time. Therefore, the optimal deletion problem is

$$\min_{l \in \{1, \dots, N\}} p_l \min_{j \neq l} c(\omega_l, \omega_j) \quad (4.17)$$

If the minimum is attained at  $l_* \in \{1, \dots, N\}$ , i.e., the scenario  $\omega_{l_*}$  is deleted, the optimal redistribution rule is  $\bar{q}_l = p_l$  for each  $l \notin \{l_*, j(l_*)\}$  and  $\bar{q}_{j(l_*)} = p_{j(l_*)} + p_{l_*}$ , where  $j(l_*) \in \arg \min_{j \neq l_*} c(\omega_{l_*}, \omega_j)$ . The optimal deletion of a single scenario will be repeated recursively until a prescribed number  $k$  of scenarios is deleted.

## **CHAPTER 5      CASE STUDY**

A numerical experiment applying the proposed methodology to an example problem of a road network is carried out in the case study. The characteristics of the test problem and some implementation details are discussed. The benefit of using the stochastic programming approach over a deterministic approach is highlighted. The computational result is commented and the proposed algorithm is examined in terms of trade-offs between computational effectiveness and solution quality. Test runs were programmed in MATLAB and performed on a standard desktop computer with 1 GB of memory and a 3.4 GHz CPU.

### **5.1 CASE STUDY DATA SET**

The road network in Dallas District is used for the case study with data taken from the Pavement Management Information System (PMIS) developed and maintained by TxDOT. The PMIS is an automated system for storing, retrieving, analyzing, and reporting pavement condition information. It can be used to retrieve and analyze pavement information to compare maintenance and rehabilitation treatment alternatives, monitor current pavement conditions, and estimate total pavement needs. The main characteristics of the Dallas District road network are presented as follows.

### 5.1.1 Size of the Network

In the PMIS database, the road network of Dallas District has five different functional class highways: Business Road (BR), Farm to the Market (FM), Interstate Highway (IH), State Highway (SH) and US Highway (US). According to their similarities in terms of the deterioration pattern, the network is divided into three broader categories as presented in Table 5.1.

**Table 5.1 Road Network Length**

<b>Highway Groups</b>	<b>Length (Lane-Kilometers)</b>
Group I (IH, US and BR)	8299
Group II (SH)	3104
Group III (FM)	5045

### 5.1.2 Planning Horizon

The objective of the case study is to develop a five-year maintenance plan for the road network, where the maintenance treatments will be applied at the beginning of each year.

### 5.1.3 Performance Indicator

In this case study, the Condition Score (CS) in the PMIS database is used as the performance indicator. The TxDOT PMIS stores three scores that represent the general condition of a pavement (TxDOT 2000). The Distress Score (DS) reflects the amount of visible surface deterioration of a pavement, with a range from 1 (the most distress) to 100 (the least distress). The Ride Score (RS) is a measure of the pavement's roughness, ranging from 0.1 (the roughest) to 5.0 (the smoothest). The Condition Score represents the pavement's overall condition in terms of both distress and ride quality ranging from 1 (the worst condition) to 100 (the best condition). The condition of a pavement is discretized into five different states according to its condition score (Table 5.2)

**Table 5.2 PMIS Condition Scores**

<b>Condition Score</b>	<b>Description</b>
90-100	Very Good
70-89	Good
50-69	Fair
35-49	Poor
1-34	Very Poor

The initial condition of the road network in terms of the percentage in each condition state is shown in Table 5.3. The numbers in this table represent the

percentage of the corresponding road type in a specific condition state. For example, 73 percent of Type I road pavements—which comprise the majority of the road network—are in “Very Good” condition.

**Table 5.3 Road Network Initial Condition (%)**

<b>Condition State\Road Groups</b>	<b>IH, US and BR</b>	<b>SH</b>	<b>FM</b>
Very Good	73	58	62
Good	11	15	16
Fair	7	10	10
Poor	5	9	8
Very Poor	4	8	4

The goal of the road network’s five year maintenance plan is that 90 percent of the road group I should be in “Very Good” condition state, and 80 percent of road groups II and III should be in “Very Good” condition state as shown in Table 5.4.

**Table 5.4 Road Condition Requirements**

<b>Condition State\Road Groups</b>	<b>I</b>	<b>II</b>	<b>III</b>
Very Good (100-80)	90%	80%	80%

**5.1.4 Transition Probability**

Generally, there are two ways that transition probability can be estimated. The first way is by simulation through pavement design equations (Gao and Zhang 2007), while the second way is by estimating the probability using historical data (Butt et al. 1987; Jiang et al. 1989; Wang et al. 1994). In this case study, the transition probability for each of the road groups is calculated based on the historical data from the Dallas PMIS database. The results are shown in Table 5.5, Table 5.6, Table 5.7.

**Table 5.5 TPM for Road Group I**

<b>Initial State\Next State</b>	<b>Very Good</b>	<b>Good</b>	<b>Fair</b>	<b>Poor</b>	<b>Very Poor</b>
Very Good	0.85	0.10	0.03	0.01	0.00
Good	0.00	0.57	0.28	0.12	0.04
Fair	0.00	0.00	0.47	0.39	0.13
Poor	0.00	0.00	0.00	0.56	0.44
Very Poor	0.00	0.00	0.00	0.00	1.00



**Table 5.6 TPM for Road Group II**

<b>Initial State\Next State</b>	<b>Very Good</b>	<b>Good</b>	<b>Fair</b>	<b>Poor</b>	<b>Very Poor</b>
Very Good	0.74	0.16	0.07	0.03	0.01
Good	0.00	0.35	0.37	0.21	0.07
Fair	0.00	0.00	0.45	0.44	0.11
Poor	0.00	0.00	0.00	0.55	0.45
Very Poor	0.00	0.00	0.00	0.00	1.00

**Table 5.7 TPM for Road Group III**

<b>Initial State\Next State</b>	<b>Very Good</b>	<b>Good</b>	<b>Fair</b>	<b>Poor</b>	<b>Very Poor</b>
Very Good	0.77	0.14	0.06	0.03	0.01
Good	0.00	0.36	0.39	0.19	0.06
Fair	0.00	0.00	0.38	0.43	0.19
Poor	0.00	0.00	0.00	0.41	0.59
Very Poor	0.00	0.00	0.00	0.00	1.00

### **5.1.5 Maintenance Effect**

Maintenance treatments could be from the least expensive in preventive maintenance to the most expensive in reconstruction. However, it is not necessary for programming at the network level to be as detailed as at the project level. Four maintenance treatments levels are used in this case study: Do Nothing, Preventive Maintenance, Light Rehabilitation and Heavy Rehabilitation. For a given section at any given year, four possible treatments can be performed. Preventive maintenance, including seal coat, micro-surfacing or thin overlay, is aimed at extending the life of bituminous surfaces by retarding the effects of weathering and aging before significant amounts of distress have occurred. Rehabilitation involves heavier treatments intended to increase the structural capacity, restore ride and seal the base and subgrade layers. For demonstration purpose, the assumed maintenance treatments effect for a pavement section is given in Table 5.8.

**Table 5.8 Maintenance Treatments Effect**

<b>M&amp;R treatment</b>	<b>Condition state before treatments</b>	<b>Condition state after treatment</b>
Do Nothing	Very Good	Very Good
	Good	Good
	Fair	Fair
	Poor	Poor
	Very Poor	Very Poor
Preventive Maintenance	Very Good	Very Good
	Good	Very Good
	Fair	Fair
	Poor	Poor
	Very Poor	Very Poor
Light Rehabilitation	Very Good	Very Good
	Good	Very Good
	Fair	Very Good
	Poor	Good
	Very Poor	Fair
Heavy Rehabilitation	Very Good	Very Good
	Good	Very Good
	Fair	Very Good
	Poor	Very Good
	Very Poor	Very Good

### 5.1.6 Maintenance Cost (Agency Cost)

The unit costs for all types of treatments are taken from the work of Wang et al. (2003) as listed in Table 5.9.

**Table 5.9 Maintenance Treatment Costs**

<b>Road Group</b>	<b>Maintenance Treatment</b>	<b>Maintenance treatment unit cost (\$1000/lane/km)</b>
I	Do Nothing	0
	Preventive Maintenance	10
	Light Rehabilitation	100
	Heavy Rehabilitation	500
II	Do Nothing	0
	Preventive Maintenance	8
	Light Rehabilitation	80
	Heavy Rehabilitation	400
III	Do Nothing	0
	Preventive Maintenance	5
	Light Rehabilitation	20
	Heavy Rehabilitation	100

## 5.2 SOLUTION OF ALD

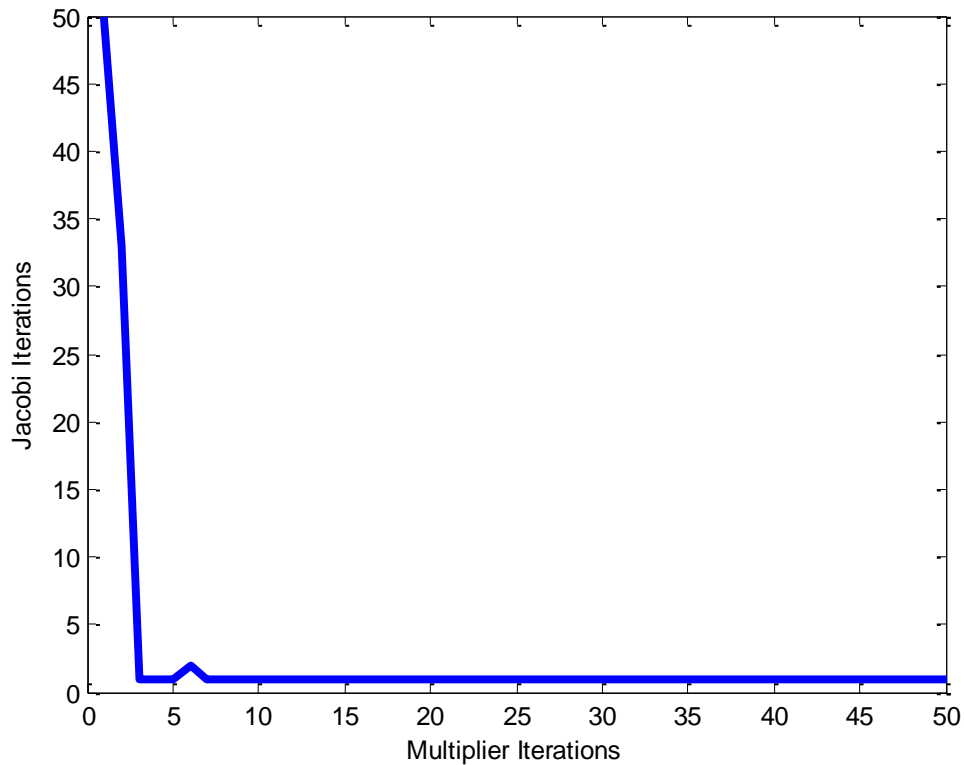
In this research, it is assumed that the budget at every year is unknown but will be allocated from \$80 million, \$100 million, and \$120 million with equal probabilities. Therefore, there are a total of 243 scenarios for this problem. The characteristics of the original problem and the subproblems after decomposition are summarized in Table 5.10. Before decomposition, the stochastic programming problem has 287,955 constraints and 24,300 variables, making it impossible to be solved on a standard desktop computer. The subproblems require much less computational effort with each subproblem having only 75 constraints and 300 variables. As noted, by adopting the decomposition technique, the size of the problem is dramatically reduced.

**Table 5.10 Computational Characteristics**

	<b>Subproblem after applying augmented Lagrangian decomposition</b>	<b>Original problem</b>
Number of constraints	75	287,955
Number of variables	300	24,300

A stopping criterion  $\varepsilon = 10^{-3}$  is used for both the method of multipliers and the Jacobi method. The value of  $\rho$  and  $\tau$  is set at 0.5. The initial values of the

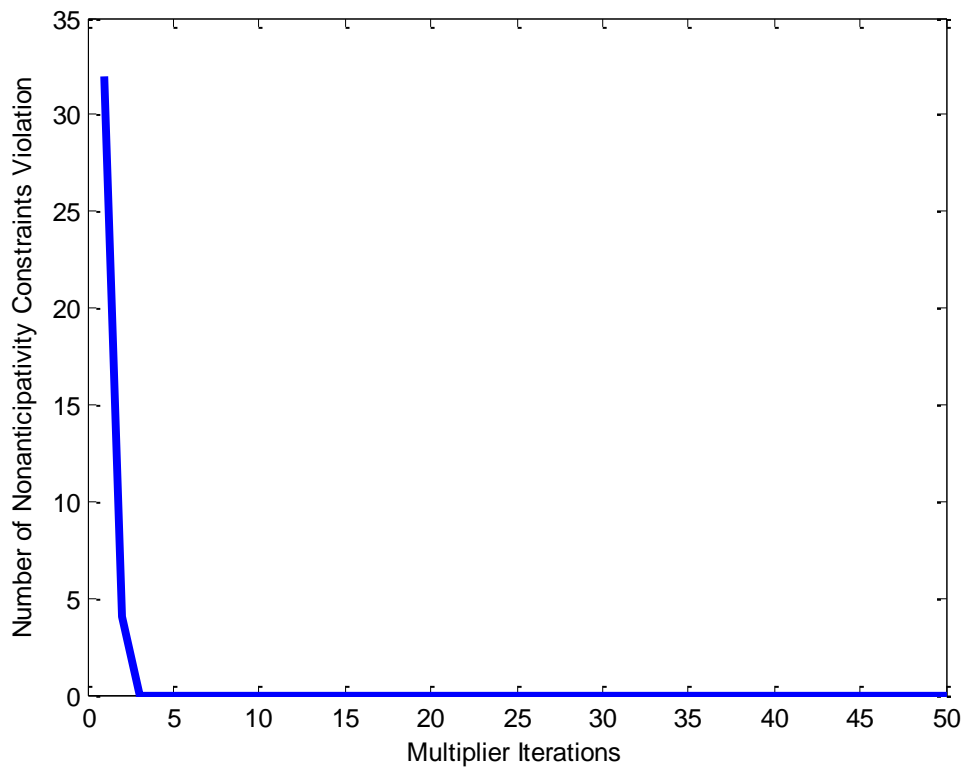
decision variable  $x^0$  are assigned zeros. Figure 5-1 shows the relationship between the iteration of the method of multipliers and the iteration of the Jacobi method. As illustrated, the Jacobi method occurs with greatest frequency at the beginning of the algorithm, then the iteration of Jacobi steps decreases rapidly.



**Figure 5-1 Number of Jacobi Steps in Each Outer Loop**

Figure 5-2 shows the relationship between the multiplier iterations and the number of nonanticipativity constraints violated. As seen in this figure, the constraints violation drops quickly during the first four multiplier iterations; then it is subsequently followed by a slower convergence until the stopping criterion is

reached. The optimal objective function value of the stochastic programming approach is obtained as 78.23.



**Figure 5-2 Number of Violated Nonanticipativity Constraints**

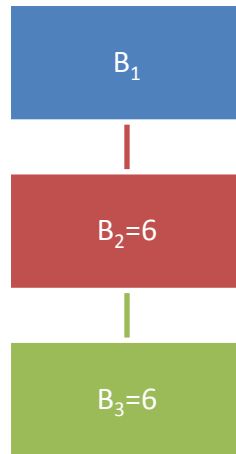
### **5.3 DETERMINISTIC SOLUTION (EV)**

An alternative to the stochastic programming (SP) approach is to consider only the expected budget values, which is known as the expected value (EV) approach. This approach is to schedule the maintenance activities assuming that

the budget will take their expected values during the planning horizon. The EV approach can be mathematically expressed as

$$EV = \min_x z(x, \bar{\xi}) \quad (5.1)$$

where  $x$  represents the decision variables,  $z$  represents the objective function and  $\bar{\xi}$  is the expected value of the random variable  $\xi$ . Using the example of Figure 3-1, the concept of the EV approach can be illustrated in Figure 5-3.



**Figure 5-3 Scenario Tree of EV Approach**

The advantage of this approach is that it is computationally easy to solve. In this research, by solving the deterministic problem (3.4)-(3.8), the detailed maintenance plan is obtained as shown in Table 5.11. The numbers in the table



are the percentage of whole road network that will receive the corresponding maintenance treatments.

**Table 5.11 Maintenance Plan of Deterministic Solution**

<b>Year</b>	<b>Do Nothing</b>	<b>Preventive Maintenance</b>	<b>Light Rehabilitation</b>	<b>Heavy Rehabilitation</b>
1	0.257	0.670	0.073	0.000
2	0.107	0.869	0.017	0.008
3	0.100	0.893	0.000	0.007
4	0.084	0.880	0.029	0.007
5	0.142	0.815	0.038	0.005

The objective function of the EV approach is obtained as 89.97, which is much better than the SP solution 78.23. This is no surprise, since the EV approach only considers one scenario while the SP considers all 243 scenarios. The EV result actually represents the upper bound of the SP problem. However, ignoring the random characteristics of future budget may lead to suboptimal result. The EV solution is infeasible (in terms of budget constraint satisfaction) to some of scenarios. As a result, some of the planned maintenance activities may have to be canceled and a new maintenance plan has to be made. In order to evaluate the benefit of using the SP method against the EV approach quantitatively, the EV solution  $\bar{x}(\bar{\xi})$  is used to calculate the expected objective

function value for all possible scenarios. The resulted quantity is called expected result of using the EV solution (EEV).

$$EEV = E_{\xi} \left( z \left( \bar{x} \left( \bar{\xi} \right), \xi \right) \right) \quad (5.2)$$

EEV measures how  $\bar{x}(\bar{\xi})$  performs, allowing subsequent-stages decisions to be chosen optimally. In other words, EEV represents the expected objective function value if decisions are made ignoring the budget uncertainty. By using (5.2), the EEV of the test problem can be calculated as 67.35. The difference between the EEV and the SP solution is called value of the stochastic solution (VSS),

$$VSS = EEV - SP \quad (5.3)$$

A small VSS means that the approximation of the SP by the EV approach is applicable. For the test example, however, VSS is almost 15% of the value of SP, which confirms that there is an obvious benefit in using a stochastic model than a deterministic one.

In order to identify the difference between SP and EV, Table 5.12 compares the maintenance plans of them at the first year. As can be seen in Table 5.12, more resources are allocated to preventive maintenance in the stochastic programming approach. Therefore, the stochastic solution alleviates the effect of possible funding shortages by allocating more resources to

preventive maintenance treatments. The underlying strategy of the stochastic solution is to spread out current funding among more pavement sections given the existence of budget uncertainty in future years. Using this strategy, the expected condition of a road network can be optimized. In practice, the proposed stochastic programming problem must be solved every year when decision makers become aware of specific appropriations and budget constraints. The maintenance plan obtained for the first year can be used to schedule activities during the year under consideration.

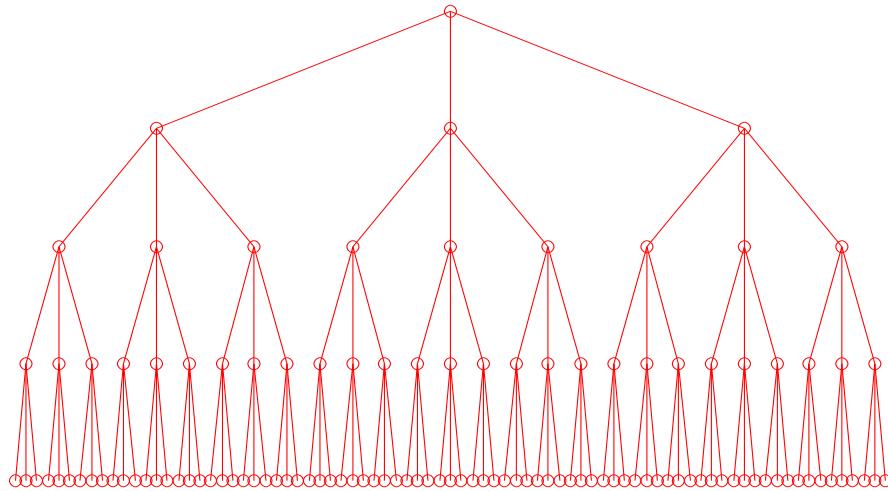
**Table 5.12 Maintenance Plan Comparison of Year 1 between EV and SP**

<b>Method</b>	<b>Do Nothing</b>	<b>Preventive Maintenance</b>	<b>Light Rehabilitation</b>	<b>Heavy Rehabilitation</b>
EV	0.257	0.670	0.073	0.000
SP	0.170	0.780	0.050	0.000

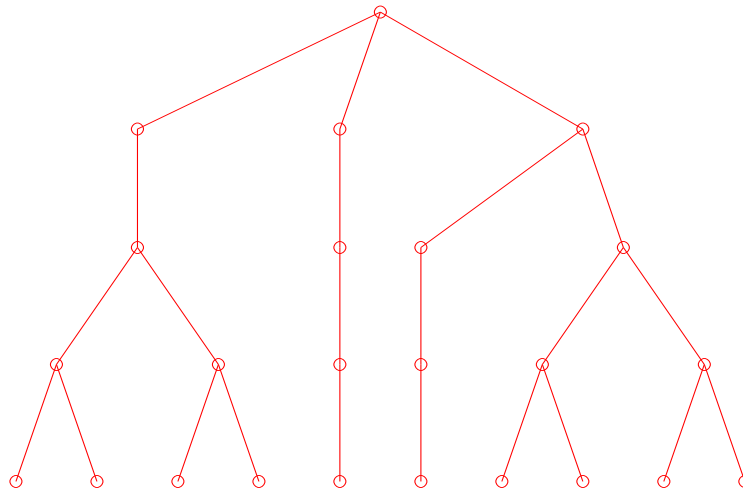
#### **5.4 COMPUTATIONAL COMPARISON (ALD, EV AND SR)**

To demonstrate the effectiveness of the proposed decomposition algorithm, a computational comparison is carried out. Table 5.13 lists the computational characteristics of the augmented Lagrangian decomposition (ALD) method, scenario reduction (SR) method and the EV approach for problems (3.11)–(3.16). The scenario reduction method is another approach to solve the stochastic programming problem. It generates a scenario subset closest to the initial

distribution in terms of a natural probability metric. In other words, only a portion of the original scenarios are selected in SR to reduce the size of the problem. As can be seen in Figure 5-4 and Figure 5-5, the size of the problem can be largely reduced after applying the scenario reduction method.



**Figure 5-4 Scenario Tree before Reduction**



**Figure 5-5 Scenario Tree after Reduction**

As shown in Table 5.13, the EV approach and the SR approach are much faster than the ALD method in terms of computational time. Because of the reduction of uncertainty, the objective function values of SR and EV are higher than the result of ALD. However, as shown in the fourth column of Table 5.13, by using the idea of (5.2), the ALD approach produces the best expected objective function value for all 243 scenarios. This is because the ALD approach takes all scenarios into consideration at the beginning of the planning horizon; and the solution of ALD consists of maintenance plan for every scenario. However, the solutions obtained from SR and EV considers only part of the scenarios. As a result, some of the planned maintenance activities may have to be re-planned in the future, which makes the solution suboptimal. It is up to the decision maker's choice to make trade-offs of solution quality and computational effort.

**Table 5.13 Computational Characteristics Comparison**

<b>Methods</b>	<b>Computational Time (seconds)</b>	<b>Objective Function Value</b>	<b>Expected Objective Function value</b>
ALD	240.1398	78.23	78.23
SR(reduced to 10 scenarios)	7.6875	80.66	72.49
EV	0.5938	89.97	67.35

## **CHAPTER 6 CONCLUSIONS AND RECOMMENDATIONS**

### **6.1 SUMMARY**

The main objective of this study is to define a methodological framework for infrastructure asset management maintenance scheduling problem under budget uncertainty and to develop solution algorithms to solve the proposed problem. A multistage linear stochastic programming model is developed and the effectiveness and efficiency of three different solution approaches are investigated. Finally, the applicability of the developed model and solution algorithms are demonstrated with solving some practical example problems.

### **6.2 CONCLUSIONS**

Conclusions drawn from this study are as follows:

1. Stochastic programming methods can be used to model the uncertainty of future maintenance budgets as random variables in infrastructure maintenance scheduling problems for better resource allocation. Stochastic programming is based on probability theory and mathematical programming. A Stochastic Programming problem can be translated to a deterministic optimization problem by defining a scenario tree. However, as the number of planning stages and number of scenarios at each stage increase, the size of the resulting deterministic problem increases quickly. Three different approaches (Augmented Lagrangian Decomposition (ALD),

Expected Value (EV) and Scenario Reduction (SR)) are investigated in this research. The ALD approach is able to produce the best results.

2. A road network example is studied as part of this research. The findings show that the proposed model and solution procedure is able to solve the maintenance scheduling problem efficiently and effectively. The benefit of using the stochastic programming approach over a deterministic approach is also discussed. Stochastic programming solutions, which take future budget uncertainty into consideration, tend to allocate more resource into preventive maintenance than deterministic solution that ignores the uncertainty information. The proposed methodology can help decision makers effectively obtain optimal maintenance planning under budget uncertainty.

### **6.3 RECOMMENDATIONS FOR FUTURE RESEARCH**

In the following, some areas are given with respect to opportunities for future research.

#### **6.3.1 Stochastic Integer Programming**

The current framework is based on stochastic linear programming, where the decision variables determine the percentage of infrastructure system receiving a certain type of maintenance treatment. As discussed in section 2.2.2, this formulation has its advantage that the solution is guaranteed to be global optimal.



However, this approach simplifies the decision making process by giving maintenance plans for “groups” instead of individual facilities. Therefore, an agency that manages an infrastructure system has to further allocate resource from “groups” to specific facilities after running the linear programming model. In other words, the Integer Programming model, whose solutions directly specify the location, timing and treatment type, may produce a better plan than the Linear Programming approach. As discussed before, the disadvantage of the IP approach is that the size of the problem increase exponentially as the number of facilities, the number of planning stages and the number of maintenance treatments increase. It is of great value to develop solution algorithms that can solve large-scale IP models for infrastructure maintenance scheduling, especially with applications under the stochastic settings.

### **6.3.2 Uncertainties other than Budget**

There are other uncertainties associated with data, models, and processes in the infrastructure asset management. For example, infrastructure deterioration is a dynamic, complicated, and stochastic process affected by a variety of factors such as usage, environmental conditions, and structural capacities, as well as certain unobserved factors. Hence, the performance of an infrastructure facility can never be predicted with absolute certainty. Ignoring such uncertainties during the modeling process may compromise the validity of an optimal solution. It is also important to take those uncertainties into consideration when making maintenance scheduling decisions.

### **6.3.3 Different Ownership**

In the current framework, the developed model is suitable for government agencies like state DOTs. In recent years, public private partnership (PPP) is becoming an increasingly popular method of funding large infrastructure projects. These PPP projects involve financing for different stages of a project including the design, build, expansion, upgrade and operation. This relatively new type of mechanisms for funding infrastructure projects has highlighted some of the challenges and issues when planning maintenance activities. Therefore, it is important that this new change being reflected in the maintenance scheduling model.

### **6.3.4 Balance between different regions**

In this dissertation, the developed model can help decision makers allocate funds to infrastructure facilities under their jurisdiction. However, for some agencies, balancing resources between different regions or districts is a practical issue that has to be taken into consideration. For example, in Texas Department of Transportation, funds have to be distributed to 25 districts and the districts can further allocate it to specific projects. Therefore, it is important to incorporate this information as additional constraints to the model, allowing maintenance plans to be developed with the consideration of different conditions and demands among districts.

### **6.3.5 Multiple objectives**

In the current methodology framework, only one objective is considered in the optimization problem formulation. Single-objective optimization is adequate if the decision maker is satisfied with optimizing only one objective. In practice, there may be more than one objective that needs to be optimized in the infrastructure maintenance scheduling process. Different competing objectives may have significantly different impacts on the resulting solutions. For example, an agency may wish to find maintenance strategies that maximize system conditions while also minimizing the maintenance expenditures. A trade-off compromise can be used to either optimize one objective and include the competing objectives as constraints, or optimize the sum of the competing objectives. In future research works, these multi-objective issues should be addressed.

## REFERENCES

- Abaza, K.A., Ashur, S.A., and Al-Khatib, I.A. (2004), "Integrated Pavement Management System with a Markovian Prediction Model", *Journal of Transportation Engineering*, 130(1), pp. 24–33.
- American Association of State Highway Officials (AASHO), (1962), *The AASHO road test report 5, pavement research. Special Report 61*, National Research Council, Washington, D.C.
- Birge, J.R. and Louveaux, F. (1997), *Introduction to Stochastic Programming*, Springer, New York.
- Blaquiere, A. (1977), *Differential Games with Piecewise Continuous Trajectories, Differential Games and Applications*, Springer-Verlag, Berlin, Germany, pp. 34–69.
- Butt, A.A. Shahin, M.Y., Feighan, K.J. and Carpenter, S.H. (1987), *Pavement Performance Prediction Model Using the Markov Process. Transportation Research Record: Journal of the Transportation Research Board*, 1123, pp. 12-19.
- Buttler, H.J. (1978), "Investment Planning of a Road Link", *Transportation Research*, Vol.12, pp. 357–366.
- Chan, W.T., Fwa, T. F. and Tan, C. Y. (1994), "Road-Maintenance Planning Using Genetic Algorithms I: Formulation", *Journal of Transportation Engineering*, 120(5), pp. 683–709.
- Dahl, G. and Minken, H. (2008), *Methods based on discrete optimization for finding road network rehabilitation strategies*, *Computers and Operations Research*, 35(7), pp. 2193–2208.

- Dupacova, J., Growe-Kuska, N. and Romisch, W. (2003), Scenario reduction in stochastic programming: an approach using probability metrics, *Mathematical Programming*, 95(3), pp. 493–511.
- Durango, P.L. and Madanat, S.M. (2002), “Optimal maintenance and repair policies in infrastructure management under uncertain facility deterioration rates: an adaptive control approach”, *Transportation Research Part A*, 36(9), pp. 763–778.
- Federal Highway Administration (FHWA), (1979), Recording and coding guide for structure inventory and appraisal of the nation’s bridges, U.S. Dept. of Transportation, Washington, D.C.
- Feighan, K.J., Shahin, M.Y. and Sinha, K. C. (1987), “A dynamic programming approach to optimization for pavement management systems”, 2nd North Am. Conf. on Managing Pavements. pp. 2.195–2.206.
- Feighan, K. J., Shahin, M.Y., Sinha, K.C. and White, T.D. (1988), “Application of Dynamic Programming and Other Mathematical Techniques to Pavement Management Systems”, *Transportation Research Record 1200*, Transportation Research Board, Washington, D.C., pp. 90–98.
- Fernandez, J.E. and Friesz T.L. (1981), “Influence of Demand-Quality Interrelationships on Optimal Policies for Stage Construction of Transportation Facilities”, *Transportation Science*, 15(1), pp. 16–31.
- Ferreira, A., Antunes, A. and Picado-Santos, L. (2002), “Probabilistic Segment-Linked Pavement Management Optimization Model”, *Journal of Transportation Engineering*, 128(6), pp. 568–577.
- Friesz, T.L. and Fernandez, J.E. (1979), “A Model of Optimal Transport Maintenance with Demand Responsiveness”, *Transportation Research Part B*, 13(4), pp. 317–339.

- Fwa, T.F., Sinha, K.C. and Riverson, J.D.N. (1988), "Highway routine maintenance programming at network level", *Journal of Transportation Engineering*, 114(5), pp. 539–554.
- Gao, L., Xie, C. and Zhang, Z. (2010), *Network-Level Multi-Objective Optimal Maintenance and Rehabilitation Scheduling*, CD Proceedings of the Transportation Research Board's 89th Annual Meeting, Washington, D.C.
- Gao, L. and Zhang, Z. (2007), "Using Markov Process and Method of Moments for Optimizing Management Strategies of Pavement Infrastructure", CD Proceedings of the Transportation Research Board 86<sup>th</sup> Annual Meeting, Washington, D.C.
- Gao, L. and Zhang, Z. (2008), "Robust Optimization for Managing Pavement Maintenance and Rehabilitation", *Transportation Research Record 2084*, Transportation Research Board, Washington, D.C., pp. 55–61.
- Gao, L. and Zhang, Z. (2008), "An Approximate Dynamic Programming Approach to Network-Level Budget Planning and Allocation for Pavement Infrastructure", CD proceeding of the Transportation Research Board 88<sup>th</sup> Annual Meeting, Washington, D.C.
- Gao, L. and Zhang, Z. (2010), "The Effect of Budget Uncertainty on Road Network Condition Fluctuation", CD proceeding of the Transportation Research Boards 89<sup>th</sup> Annual Meeting, Washington, D.C.
- Gao, L., Aguiar-Moya, J.P., and Zhang, Z. (2011), "A Bayesian Analysis of Heterogeneity in Modeling of Pavement Fatigue Cracking", accepted for publication at *Journal of Computing in Civil Engineering*.
- Gao, L., Aguiar-Moya, J.P., and Zhang, Z. (2011), "Performance Modeling of Infrastructure Condition Data with Maintenance Intervention", accepted for publication at *Journal of the Transportation Research Record*.

- Garcia, D.A. and Riggins, M. (1984), Serviceability and Distress Methodology for predicting pavement performance. *Transportation Research Record: Journal of the Transportation Research Board*, 997, pp.56-61.
- Golabi, K., Kulkarni, R. B., and Way, G. B. (1982), A statewide pavement management system, *Interfaces*, 12(1), pp.5–21.
- Greene, W.H. (2007), *Econometric Analysis* (6th ed.), Macmillan, New York.
- Guignier, F. and Madanat, S.M. (1999), Optimisation of infrastructure systems maintenance and improvement policies, *Journal of Infrastructure Systems*, 5(4).
- Guignier, F. and Madanat, S.M. (2000), “Optimization of Infrastructure Systems Maintenance and Improvement Policies”, *Transportation Research Circular 498*, Transportation Research Board, Washington, D.C., pp. 124–134.
- Guo, R. and Gao, L. (2011), “Optimal Pavement Maintenance Policy under Warranty Contract”, *CD Proceedings of The Transportation Research Boards’ 90th Annual Meeting*, Washington, D.C.
- Hong and Prozzi (2010), Roughness model accounting for heterogeneity based on in-service pavement performance data, *Journal of Transportation Engineering*, 136(3), pp. 205–213.
- Heitsch, H. and Romisch, W. (2009), Scenario tree modeling for multistage stochastic programs, *Mathematical Programming*, 118(2), pp. 371–406.
- Humplick, F. (1992), Highway Pavement Distress Evaluation: Modeling Measurement Error, *Transportation Research Part B*. 26(2), pp.135-154.

- Jacobs, T.L. (1992), "Optimal long term scheduling of bridge deck replacement and rehabilitation", *Journal of Transportation Engineering*, 118(2), pp. 312–322.
- Jiang, Y., Saito, M. and Sinha, K.C. (1989), "Bridge Performance Prediction Model Using the Markov Chain", *Transportation Research Record* 1180, Transportation Research Board, Washington D.C., pp. 25–32.
- Jido, M., Otazawa, T. and Kobayashi, K. (2008), "Optimal Repair and Inspection Rules Under Uncertainty", *Journal of Infrastructure Systems*, 14(2), pp. 150–158.
- Karabakal, N., Bean, J. C. and Lohmann, J. R. (1994), *Scheduling Pavement Maintenance with Deterministic Deterioration and Budget Constraints*, Technical Report 94-18, University of Michigan.
- Li, N., Xie, W. and Haas, R. (1996), "Reliability-Based Processing of Markov Chains for Modeling Pavement Network Deterioration", *Transportation Research Record* 1524, Washington, D.C., pp. 203–213.
- Li, Y. and Madanat, S.M. (2002), "A steady-state solution for the optimal pavement resurfacing problem", *Transportation Research Part A*, 36(6), pp. 525–535.
- Li, Z. and Puyan, M. (2006), "A stochastic optimization model for highway project selection and programming under budget uncertainty", *Proceeding of the Applications of Advanced Technology in Transportation*.
- Liu, C., Hammad, A. and Itoh, Y. (1997), "Multiobjective optimization of Bridge Deck Rehabilitation Using a Genetic Algorithm", *Microcomputers in Civil Engineering*, 12(6), pp. 431–443.
- Madanat, S.M., Bulusu, S. and Mahmoud, A. (1995), "Estimation of Infrastructure Distress Initiation and Progression Models", *Journal of Infrastructure Systems*, 1(3), pp.247-256.



- Madanat, S.M., Karlaftis, M.G. and McCarthy, P.S. (1997), Probabilistic Infrastructure Deterioration Models with Panel Data, *Journal of Infrastructure Systems*. 3(1), pp. 4-9. 1997.
- Madanat, S.M., Park, S. and Kuhn, K. (2006), "Adaptive Optimization and Systematic Probing of Infrastructure System Maintenance Policies under Model Uncertainty", *Journal of Infrastructure Systems*, 12(3), pp. 192–198.
- Markow, M. J., Brademeyer, B. D., Sherwood, J. and Kenis W. J. (1987), "The economic optimization of pavement maintenance and rehabilitation policy", 2nd North Am. Conf. on Managing Pavements, pp. 2.169–2.182.
- Mauch, M. and Madanat, S. (2001), Semiparametric Hazard Rate Models of Reinforced Concrete Bridge Deck Deterioration, *Journal of Infrastructure System*, 7(2), pp.49-57.
- Mishalani, R.G. and Madanat, S.M. (2002), Computation of Infrastructure Transition Probabilities Using Stochastic Duration Models. *Journal of Infrastructure Systems*, 8(4), pp.139-148.
- Ophem, H.V. and Schram, A. (1997), Sequential and Multinomial Logit: A Nested Model. *Empirical Economics*, 22, pp. 131-152.
- Ouyang, Y. and Madanat, S.M. (2004), "Optimal scheduling of rehabilitation activities for multiple pavement facilities: exact and approximate solutions", *Transportation Research Part A*, 38(5), pp. 347–365.
- Ouyang, Y. and Madanat, S.M. (2006), "An analytical solution for the finite-horizon pavement resurfacing planning problem", *Transportation Research Part B*, 40(9), pp. 767–778.
- Paterson, William D.O. (1987), *Road Deterioration and Maintenance Effects: Models for Planning and Management*. The Highway Design and Maintenance Series, The John Hopkins Univ. Press, Baltimore.

Peterson, D. E. (1985), Life-Cycle Cost Analysis of Pavements, NCHRP Synthesis 122, Transportation Research Board, Washington, D.C.

Peterson, D. E. (1987), Pavement Management Practices, NCHRP Synthesis 135, Transportation Research Board, Washington, D.C.

Pilson, C., Hudson, W. R. and Anderson, V. (1999), "Multiobjective Optimization in Pavement Management by Using Genetic Algorithms and Efficient Surfaces", Transportation Research Record 1655, Transportation Research Board, Washington, D.C., pp. 42–48.

Prozzi, J.M. (2001), Modeling Pavement Performance by Combining Field and Experimental Data, Dissertation, The University of California at Berkley.

Rempala, R. and Zabczyk, J. (1988), On the maximum principle for deterministic impulse control problems, Journal of Optimization Theory and Applications, 59(2), pp. 281-288.

Robelin, C.A. and Madanat, S.M. (2006), Dynamic Programming based Maintenance and Replacement Optimization for Bridge Decks using History-Dependent Deterioration Models, AATT.

Rosa, C.H. and Ruszczyński, A. (1996), on augmented Lagrangian decomposition methods for multistage stochastic programs, Annals of Operations Research, 64(2), pp. 289–309.

Ruszczyński, A. (1997), decomposition methods in stochastic programming, Mathematical Programming, 79(1), pp. 333–353.

Smilowitz, K. and Madanat, S.M. (2000), "Optimal Inspection and Maintenance Policies for Infrastructure Networks", Computer-Aided and Infrastructure Engineering, 15(1), pp. 5–13.

- Stampley, B. E., Miller, B., Smith, R.E. and Scullion T. (1995), Pavement Management Information System Concepts, Equations, and Analysis Models, TxDOT Research Report 1989-1, Texas Transportation Institute, Texas A&M University System.
- Sun, W. and Yuan, Y. (2006), Optimization Theory And Methods Nonlinear Programming, Springer, New York.
- Tsunokawa, K. and Schofer, J.L. (1994), "Trend Curve Optimal Control Model For Highway Pavement Maintenance: Case Study and Evaluation", Transportation Research Part A, 28(2), pp. 151–166.
- Texas 2030 Committee (2011), It's About Time: Investing in Transportation to Keep Texas Economically Competitive.
- TxDOT (2000), Managing Texas Pavements, Design Division of the Texas Department of Transportation, Austin, Texas.
- Wang, K.C.P., Zaniewski, J. and Way, G. (1994), "Probabilistic behavior of pavements", Journal of Transportation Engineering, 120(3), pp. 385–375.
- Wang, F., Zhang, Z. and Machemehl, R.B. (2003), "Decision-Making Problem for Managing Pavement Maintenance and Rehabilitation Projects", Transportation Research Record 1853, Transportation Research Board, Washington, D.C., pp. 21–28.
- Wang Y.H., Mahboub, K.C. and Hancher, D.E. (2005), survival analysis of fatigue cracking for flexible pavements based on long-term pavement performance data, Journal of Transportation Engineering, 131(8), pp.608-616.
- Watanatada, T., Paterson, W., Bhandari, A., Harral, C., Dhareshwar, A., and Tsunokawak, K. (1987), Volume 1 Description of the HDM III model, World Bank, Washington, D.C.

Wu, Z. and Flintsch, G.W. (2009), "Pavement Preservation Optimization Considering Multiple Objectives and Budget Variability", *Journal of Transportation Engineering*, 135(5), pp. 305–315.

Zhang, Z. and Damnjanovic, I. (2006), Applying Method of Moments to Model Reliability of Pavements Infrastructure, *Journal of Transportation Engineering*, 132(5), pp. 416-424.

Zhang, Z. and Gao, L. (2010), "Predicting Pavement Deterioration: A Sequential Logit Model Approach", CD Proceedings of the Transportation Research Boards' 89th Annual Meeting, Washington, D.C.

## VITA

Lu Gao attended Bashu Middle School in Chongqing, China. In 2001, he entered Tsinghua University in Beijing, China. He also studied at the Nanyang Technological University in Singapore as an exchange student in 2004. Lu graduated with his Bachelor of Engineering in Civil Engineering from Tsinghua University in June, 2005. Soon after, he joined the transportation engineering program at the University of Texas at Austin. He received his Masters of Science in Transportation/Civil Engineering from the University of Texas at Austin in May, 2007.

Permanent email: [gaolu@mail.utexas.edu](mailto:gaolu@mail.utexas.edu)

This manuscript was typed by the author.