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► **To cite this version:**

Laurence Cholvy, Laurent Perrussel, William Raynaut, Jean-Marc Thévenin. Consistency-Based Reliability Assessment. European Conference on Symbolic and Quantitative Approaches to Reasoning with Uncertainty (ECSQARU 2015), Jul 2015, Compiègne, France. Symbolic and Quantitative Approaches to Reasoning with Uncertainty, pp. 246-256, 2015. <hal-01316842>

HAL Id: hal-01316842

<https://hal.archives-ouvertes.fr/hal-01316842>

Submitted on 17 May 2016

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To cite this version : Cholvy, Laurence and Perrussel, Laurent and Raynaut, William and Thévenin, Jean-Marc *Consistency-Based Reliability Assessment*. (2015) In: European Conference on Symbolic and Quantitative Approaches to Reasoning with Uncertainty (ECSQARU 2015), 15 July 2015 - 17 July 2015 (Compiègne, France).

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Consistency-Based Reliability Assessment

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Abstract. This paper addresses the question of assessing the relative reliability of unknown information sources. We propose to consider a phase during which the consistency of information they report is analysed, whether it is the consistency of each single report, or the consistency of a report w.r.t. some trusted knowledge or the consistency of different reports together. We adopt an axiomatic approach by first giving postulates which characterize how the resulting reliability preorder should be; then we define a family of operators for building this preorder and demonstrate that it satisfies the proposed postulates.

1 Motivation

Techniques for merging raw information have been studied in an extensive way. These techniques usually assume that all information provided by the sources (i.e. agents) should be considered as a whole. Two different approaches have been studied: the first one considers sources in an equal way and has led to merging techniques such as majority, arbitration merging or distance-based merging for solving conflict between contradicting information [9, 5, 6]. The second one distinguishes sources through a reliability criterion. Taking sources reliability into account provides rationales for discounting or ignoring pieces of information whose source is not considered as sufficiently reliable. Some promote a quantitative model of reliability: information sources are associated with a reliability level represented by a number used by the merging operator. According to the belief function theory, the reliability level of a source is a number between 0 and 1. This number is then used by the discounting rule in order to weaken the importance of information provided by this source [13]. Some others promote a qualitative approach to reliability and consider that information sources are ranked according to their reliability. This order or pre-order is then used by the merging operator. In [3], the author defines a merging operator which assumes that the sources are totally ordered : if s is said to be more reliable than s' and together provide contradicting information, then information provided by s is privileged; while information provided by s' which does not contradict information of s is also considered as acceptable. The same idea is followed by [10] for reasoning about more complex beliefs and in [12] for revising a belief base.

All previous works assume that the reliability of the sources is given as a parameter (quantitative or qualitative): they do not address the question of how to build up this reliability.

This paper addresses this key question. We adopt a qualitative point of view to reliability representation : the relative reliability of information sources is represented by a total preorder. We focus on the question of estimating this reliability preorder in the following context: sources are unknown (no extra information about them is available) and information provided by the sources is only qualitative (i.e., statements). We propose to consider a phase, before the information merging phase, during which information sources are observed in order to obtain a reliability preorder. We consider that during this phase, the most important is to analyse the consistency of information reported by the different sources, be it the consistency of each single report, or the consistency of a report w.r.t. some trusted knowledge, or the consistency of different reports together.

This paper is organized as follows. Section 2 presents preliminary definitions. Postulates which axiomatically characterize the reliability preorders are given in section 3. Section 4 describes a generic operator building such preorders and demonstrate that it satisfies the postulates. Examples of operators are given in section 5. Finally, section 6 concludes the paper.

2 Preliminaries

Let A be a finite set of agents; let L be a propositional logic defined over a finite set of propositional letters and propositional constants \top and \perp . An interpretation m is a mapping from the set of formulas of L to the set of truth values $\{0, 1\}$ so that $m(\top) = 1$ and $m(\perp) = 0$. The set of all interpretations is denoted M . Interpretation m is a *model* of formula F iff $m(F) = 1$. *Tautologies* are formulas which are interpreted by 1 in any interpretation. We write $\models F$ when F is a tautology. A formula is *consistent* iff it has at least one model.

Let \leq be a total preorder on A representing the relative reliability of agents: $a \leq b$ stands for b is at least as reliable as a . $GT(a, \leq) = \{x \in A \setminus \{a\} : a \leq x\}$ is the set of agents which are as least as reliable as a . Let $a \in A$, \leq_1 be a total preorder on A and \leq_2 a total preorder on $A \setminus \{a\}$; \leq_1 is *compatible with* \leq_2 iff $\forall x \forall y \ x \leq_2 y \implies x \leq_1 y$.

In the following, raw information is a *communication* consisting of a pair associating an agent and a statement:

- A *communication set on* A is a set of pairs $\langle a, \varphi \rangle$ where $a \in A$ and φ is a formula of L . $\langle a, \varphi \rangle$ means that agent a has reported φ .
- Let Ψ be a communication set on A . $Ag(\Psi) = \{a \in A, \exists \varphi \langle a, \varphi \rangle \in \Psi\}$.
- Given a communication set Ψ , we define *the communication set of* a as $\Psi_a = \{\langle a, \varphi \rangle \mid \langle a, \varphi \rangle \in \Psi\}$ and the communication of a set of agents C as

$$\Psi(C) = \bigcup_{a \in C} \Psi_a$$

The *report* associated to some Ψ represents the content of the communication:

$$\text{Report}(\Psi) = \begin{cases} \bigwedge_{\langle a, \varphi \rangle \in \Psi} \varphi & \text{if } \Psi \neq \emptyset \\ \top & \text{otherwise} \end{cases}$$

- Let Ψ and Ψ' be two communication sets on A . Ψ and Ψ' are *equivalent* iff for any agent a in A : $\models \text{Report}(\Psi_a) \leftrightarrow \text{Report}(\Psi'_a)$. That is, a 's report in Ψ is equivalent to a 's report in Ψ' . We write $\Psi \equiv \Psi'$. Obviously $\Psi \equiv \Psi'$ iff $\forall C \subseteq A \ \Psi(C) \equiv \Psi'(C)$
- Let Ψ and Ψ' be two communication sets on A . Ψ and Ψ' are *weakly equivalent* iff for any agent a in A , $\exists b, \exists c \in A$ such that $\models \text{Report}(\Psi_a) \leftrightarrow \text{Report}(\Psi'_b)$ and $\models \text{Report}(\Psi'_a) \leftrightarrow \text{Report}(\Psi_c)$. That is, we relax here the constraint that report of agent a should be equivalent both in Ψ and Ψ' ; instead we only require some other agent, possibly different from a , report equivalent information. We write $\Psi \rightleftharpoons \Psi'$. Obviously $\Psi \equiv \Psi'$ iff $\forall C \subseteq A \ \Psi(C) \rightleftharpoons \Psi'(C)$

2.1 IC-contradictory communication sets

Consistency of reports will be evaluated with respect to some integrity constraint IC , a consistent formula of L . IC has to be viewed as information taken for granted or certain. Let us now revisit the classical notion of minimal inconsistent set w.r.t. communication sets. Formally, let Ψ be a set of communications on A :

- Ψ is *IC-contradictory* iff $\text{Report}(\Psi) \wedge IC$ is inconsistent; otherwise Ψ is *IC-consistent*.
- Ψ is *minimal IC-contradictory* iff Ψ is *IC-contradictory* and no strict subset of Ψ is *IC-contradictory*.
- $\Psi \perp IC$ denotes *the set of minimal IC-contradictory subsets of Ψ* .
- $A^\perp = \cup_{F \in \Psi \perp IC} \text{Ag}(F)$ is the set of agents which have reported a piece of information which belongs to some minimal *IC-contradictory* communication set. Notice that $A^\perp \neq \emptyset$ iff Ψ is *IC-contradictory*.

Example 1. Consider agents a, b, c, d and propositional letters p, q, r, s . Assume $IC = \neg(p \wedge q)$ and $\Psi = \{\langle a, p \rangle, \langle a, r \rangle, \langle b, q \rangle, \langle c, \neg r \rangle, \langle d, s \rangle\}$. The minimal *IC-contradictory* subsets of Ψ are $E_1 = \{\langle a, p \rangle, \langle b, q \rangle\}$ and $E_2 = \{\langle c, \neg r \rangle, \langle a, r \rangle\}$. Thus $A^\perp = \{a, b, c\}$.

Example 2. Consider now agents a, b and propositional letters p, q . Assume $IC = p$ and $\Psi = \{\langle a, \neg p \wedge q \rangle, \langle b, \neg q \rangle\}$. The *IC-contradictory* subsets of Ψ are $E_1 = \{\langle a, \neg p \wedge q \rangle\}$ and $E_2 = \{\langle a, \neg p \wedge q \rangle, \langle b, \neg q \rangle\}$ but only E_1 is minimal. Thus $A^\perp = \{a\}$.

This last example shows that A^\perp is not the set of *all* agents which bring some contradiction. A^\perp is to be seen as the set of agents which prevent the consistency of Ψ i.e if communications of agents of A^\perp are ignored, Ψ becomes *IC-consistent*. I.e., $\Psi \setminus A^\perp$ is *IC-consistent* (but not necessarily maximal consistent).

2.2 IC-conflicting agents

We go further for revisiting the notion of minimal inconsistent set by considering set of agents rather than set of statements. This set helps us to identify the sources which are related to inconsistent report.

- Let $C \subseteq A$. C is *IC-conflicting* iff $\text{Report}(\Psi(C)) \wedge IC$ is inconsistent.
- C is *minimal IC-conflicting* iff it is *IC-conflicting* and no strict subset of C is *IC-conflicting*.

Example 3. Let's consider the previous example. $\{a, b, c, d\}$ is *IC-conflicting*. $\{a, b\}$ and $\{a, c\}$ are minimal *IC-conflicting*.

We can show that the union of minimal *IC-conflicting* subsets of A is included in A^\perp . But the reverse is not true: consider $IC = \neg q$ and $\Psi = \{ \langle a, p \rangle, \langle b, \neg p \rangle, \langle b, q \rangle \}$. $A^\perp = \{a, b\}$ while the only minimal *IC-conflicting* subset of agents is $\{b\}$.

3 Assessing reliability

The following postulates define in an axiomatic way that reliability assessment should be rooted in the notion of contradiction occurring in a set of communications.

Given a set of agents A , an integrity constraint IC and a communication set Ψ , the total preorder representing the relative reliability of agents in A is denoted $\Gamma^{IC,A}(\Psi)$. The operator Γ , which defines this relative reliability preorder is characterized as follows:

- P1** $\Gamma^{IC,A}(\Psi)$ is a total preorder on A .
- P2** If $\Psi \equiv \Psi'$ then $\Gamma^{IC,A}(\Psi) = \Gamma^{IC,A}(\Psi')$.
- P3** If $\models IC \leftrightarrow IC'$ then $\Gamma^{IC,A}(\Psi) = \Gamma^{IC',A}(\Psi)$.
- P4** If $\models \text{Report}(\Psi_a)$ then $\Gamma^{IC,A}(\Psi)$ is compatible with $\Gamma^{IC,A \setminus \{a\}}(\Psi \setminus \Psi_a)$.
- P5** If A is not *IC-conflicting* then $\Gamma^{IC,A}(\Psi)$ is the equality preorder.
- P6** If A is *IC-conflicting* then $A \setminus A^\perp \subseteq GT(a, \Gamma^{IC,A}(\Psi))$ for any $a \in A^\perp$.
- P7** If $\{a_1, \dots, a_k\}$ ($k \geq 2$) is a minimal *IC-conflicting* subset of agents, then $\exists i \forall j \ j \neq i, GT(a_j, \Gamma^{IC,A}(\Psi)) \subset GT(a_i, \Gamma^{IC,A}(\Psi))$.

Postulate **P1** specifies that the expected result is a total preorder. **P2** and **P3** deal with syntax independence. More precisely, if we consider two equivalent communication sets or if we consider two equivalent *IC* formulas, then we get the same total preorder on agents. **P4** states that an agent which reports a tautology or which reports no information has no influence on the relative reliability of other agents. **P5**, **P6** and **P7** focus on consistency of information provided by agents in A . Postulate **P5** considers the case when A is not *IC-conflicting* (i.e Ψ set is not *IC-contradictory*). In such a case, the sources are considered as equally reliable. **P6** and **P7** consider the cases when A is *IC-conflicting*. According to **P6**, any agent reporting a piece of information belonging to some

minimal IC -contradictory communication set is considered as less reliable than any other agent which have not. According to **P7**, if some agents are minimally IC -conflicting, then at least one of these agents is strictly less reliable than the others. This is inline with our understanding of reliability: if some agents are equally reliable, then after merging we will believe, with the same strength, information they will provide. However, it is generally assumed [4, 11] that graded belief satisfies a modal logic axiom which states that beliefs should be consistent: that is, two contradictory pieces of information cannot be believed with the same strength. Consequently, agents who are involved in a minimal IC -conflicting set cannot be equally reliable.

4 Operator assessing reliability

In this section, we propose a generic operator which builds the reliability preorder of agents by taking into account their contribution to inconsistencies.

4.1 Contribution of agents to inconsistencies

We start by introducing a measure to quantify the inconsistency degree of communication sets w.r.t. some IC . This measure is adapted from the one proposed in [8] for measuring inconsistency of sets of formulas.

Definition 1. A syntax weak-independent IC -inconsistency measure is a function I_{IC} which associates any communication set Ψ with a positive real number $I_{IC}(\Psi)$ so that:

- **Consistency** : $I_{IC}(\Psi) = 0$ iff Ψ is IC -consistent.
- **Monotony** : $I_{IC}(\Psi \cup \Psi') \geq I_{IC}(\Psi)$
- **Dominance** : for all ϕ and ψ , if $IC \wedge \phi \models \psi$ and $IC \wedge \phi$ is consistent, then $I_{IC}(\Psi \cup \{< a, \phi >\}) \geq I_{IC}(\Psi \cup \{< b, \psi >\})$ for any $a, b \in A$.
- **Free formula independence** : If $< a, \phi >$ is free (it does not belong to any minimal IC -contradictory subset of $\Psi \perp IC$), then $I_{IC}(\Psi) = I_{IC}(\Psi \setminus \{< a, \phi >\})$.
- **Syntax weak-independence** :
 1. for all IC' if $\models IC \leftrightarrow IC'$ then $I_{IC}(\Psi) = I_{IC'}(\Psi)$
 2. for all Ψ' if $\Psi \equiv \Psi'$ then $I_{IC}(\Psi) = I_{IC}(\Psi')$

The consistency property states that the measure of inconsistency of a communication set is null iff this communication set is not IC -contradictory. The monotony property says that the measure of inconsistency of a communication set does not decrease if we add more communications in this set. The dominance property states that logically stronger reports bring potentially more contradictions. The free formula independence property states that adding a report that does not cause any contradiction cannot change the consistency measure of the communication set. Finally, the syntax weak-independence says that the measure of inconsistency of a communication set does not depend on the syntax on

the integrity constraints. It also says that two weakly equivalent communication sets get the same measure of inconsistency.

Notice that $I_{IC}(\emptyset) = 0$ since $\text{Report}(\emptyset) = \top$.

Proposition 1. *Let Ψ be a communications set on A , IC an integrity constraint and $a \in A$. If $a \notin A^\perp$ then $I_{IC}(\Psi_a) = 0$. The reverse is not true.*

Next we consider a function for measuring how much an agent contributes to the IC -inconsistency of a communication set. The contribution of an agent to the fact that Ψ is IC -contradictory is defined as the Shapley value and measures the importance of this agent in a coalitional game defined by function I_{IC} [8].

Definition 2. *Consider a set of agents A , a communication set Ψ on A , an integrity constraint IC and a syntax weak-independent IC -inconsistency measure I_{IC} . Function $\text{Cont}_\Psi^{I_{IC}}$ associates any agent a with a positive real number $\text{Cont}_\Psi^{I_{IC}}(a)$ so that:*

$$\text{Cont}_\Psi^{I_{IC}}(a) = \sum_{\substack{C \subseteq A \\ C \neq \emptyset}} \frac{(|C| - 1)! (|A| - |C|)!}{|A|!} (I_{IC}(\Psi(C)) - I_{IC}(\Psi(C \setminus \{a\})))$$

Proposition 2. *Let Ψ be a communications set on A , IC an integrity constraint and $a \in A$. Then, $a \notin A^\perp \implies \text{Cont}_\Psi^{I_{IC}}(a) = 0$. The reverse is not true.*

Given the function $\text{Cont}_\Psi^{I_{IC}}$, one can obviously define a total preorder on A as follows:

Definition 3. *Let a and b be two agents of A .*

$$a \leq_{\text{Cont}_\Psi^{I_{IC}}} b \quad \text{iff} \quad \text{Cont}_\Psi^{I_{IC}}(a) \geq \text{Cont}_\Psi^{I_{IC}}(b)$$

This defines the reliability as follows: a source is considered strictly more (resp, equally) reliable than another iff its contribution to the global inconsistency is strictly smaller than (resp equal to) the contribution of the other. But, unfortunately, this preorder does not satisfy the seven postulates, as shown in the following.

Proposition 3. $\leq_{\text{Cont}_\Psi^{I_{IC}}}$ *satisfies P1–P6 but does not satisfy P7.*

To prove that $\leq_{\text{Cont}_\Psi^{I_{IC}}}$ does not satisfy **P7**, just consider $\Psi = \{ \langle a, p \rangle, \langle b, \neg p \rangle \}$ and $IC = \text{true}$. $\{a, b\}$ is a *minimal IC-conflicting* set of agents but however, $a =_{\text{Cont}_\Psi^{I_{IC}}} b$.

As a consequence, we have to find another way to build operators for reliability assessment. This is the purpose of the following paragraph.

4.2 Γ^{IC} operator

Definition 4. Consider a set of agents A , a communication set on A , Ψ , an integrity constraint IC and a given syntax weak-independent IC -inconsistency measure I_{IC} . The operator Γ^{IC} for assessing reliability is defined by:

1. $X \leftarrow A$
2. $E \leftarrow \Psi \perp IC$
3. $\leq \leftarrow \{a \leq b \mid a, b \in A\}$
4. **while** $E \neq \emptyset$ **do**
 - (a) Deterministically choose $a \in Ag(\cup_{F \in E} F)$ which maximizes $Cont_{\Psi}^{I_{IC}}(a)$
 - (b) $X \leftarrow X \setminus \{a\}$
 - (c) $E \leftarrow E \setminus \{F \in E \mid a \in Ag(F)\}$
 - (d) $\leq \leftarrow \leq \setminus \{b \leq a \mid b \in X\}$
5. **return** \leq

In the previous algorithm, X is the set of agents which has to be ordered, E is the set of minimal IC -contradictory subsets of Ψ which contain communications of agents in X . \leq is the reliability pre-order and $a \leq b$ stands for b is at least as reliable as a . First Lines 1–3 initialize the variables: X is initialized as A , E contains all the minimal IC -contradictory subsets of Ψ and \leq is equality. Then according to lines (4) and (a)–(d), the operator chooses one agent a among those which maximally contribute to the IC -contradiction of Ψ , removes a from X , deletes from E all the subsets which contains some communication of a , and updates \leq so that a is no more reliable than agents in X . This is done until E is empty. Notice that line (a) expresses a deterministic choice, such as lexicographic order.

Example 4. Consider agents a, b, c, d and propositional letters p, q, r . Consider $\Psi = \{\langle a, p \rangle, \langle b, q \rangle, \langle c, \neg q \rangle, \langle d, r \rangle\}$ and $IC = \neg r$. Consequently, $\Psi \perp IC = \{\{\langle b, q \rangle, \langle c, \neg q \rangle\}, \{\langle d, r \rangle\}\}$. Assume that:

$$Cont_{\Psi}^{I_{IC}}(a) < Cont_{\Psi}^{I_{IC}}(b) = Cont_{\Psi}^{I_{IC}}(c) < Cont_{\Psi}^{I_{IC}}(d)$$

Assume a lexicographic order for choice. First, Lines 1–3 sets X , E and \leq as follows:

$$\begin{aligned} X &= \{a, b, c, d\} & E &= \{\{\langle b, q \rangle, \langle c, \neg q \rangle\}, \{\langle d, r \rangle\}\} \\ \leq &= \{a = b = c = d\} \end{aligned}$$

Next, first iteration chooses "d" at step (a) and we get:

$$\begin{aligned} X &= \{a, b, c\} & E &= \{\{\langle b, q \rangle, \langle c, \neg q \rangle\}\} \\ \leq &= \{d < a = b = c\} \end{aligned}$$

Lexicographic order entails that the 2nd iteration chooses "b":

$$\begin{aligned} X &= \{a, c\} & E &= \emptyset \\ \leq &= \{d < b < a = c\} \end{aligned}$$

As $E = \emptyset$, the algorithm stops and returns the pre-order: $d < b < a = c$.

The following propositions show that ranks are coherent with the inconsistency measure. That is, the more an agent is inconsistent, the less it is reliable:

Proposition 4. *Let Ψ be a communications set on A and IC an integrity constraint. Let \leq be the preorder given by operator Γ_1^{IC} . For any two agents a and $b \in A$:*

$$\text{If } \text{Cont}_{\Psi}^{IC}(a) \geq \text{Cont}_{\Psi}^{IC}(b) \text{ then } a \leq b$$

We have the immediate following corollary:

Corollary 1. *If $a \notin A^\perp$ then for all $b \in A$, $b \leq a$.*

The opposite direction of previous proposition can only be considered for strict order. This is due to the choice step (a): an agent may maximize the consistency measure but may not be chosen; agents may then be considered with same rank of reliability as agents getting a lower measure related to their contribution to inconsistency. In the previous example, $a = c$ while $\text{Cont}_{\Psi}^{IC}(a) < \text{Cont}_{\Psi}^{IC}(c)$.

Proposition 5. *Let Ψ be a communications set on A and IC an integrity constraint. Let \leq be the preorder given by operator Γ_1^{IC} . For any two agents a and $b \in A$:*

$$\text{If } a < b \text{ then } \text{Cont}_{\Psi}^{IC}(a) \geq \text{Cont}_{\Psi}^{IC}(b)$$

The two previous propositions show that the choice step plays a crucial role in the behavior in the operator. Indeed this choice enforces the satisfaction of postulate **P7**.

Theorem 1. *Consider a set of agents A , a communication set on A , Ψ , an integrity constraint IC and a syntax weak-independent IC -inconsistency measure I_{IC} . Γ^{IC} operators satisfy postulates **P1-P7**.*

5 Examples of Inconsistency measures

Let us now detail two possible inconsistency measures which allow us to build two reliability assessment operators. These two measures are based on the ones proposed by [8]. The first measure is inspired by the drastic distance:

Definition 5 (Drastic measure). *Let Ψ be a set of communications on A and IC a consistent formula. The drastic inconsistency measure I_d^{IC} is defined by:*

$$I_d^{IC}(\Psi) = \begin{cases} 0 & \text{if } \Psi \text{ is } IC\text{-consistent} \\ 1 & \text{otherwise} \end{cases}$$

Theorem 2. *I_d^{IC} is a syntax weak-independent inconsistency measure.*

Then Γ_d^{IC} is a good candidate to assess reliability.

Example 5. Consider agents a, b, c, d and propositional letters p, q, r . Consider $\Psi = \{ \langle a, p \wedge q \rangle, \langle b, q \rangle, \langle c, \neg q \rangle, \langle d, r \rangle \}$ and $IC = \neg r$. Then we get: $Cont_{\Psi}^{IC}(a) = \frac{1}{12}$ $Cont_{\Psi}^{IC}(b) = \frac{1}{12}$ $Cont_{\Psi}^{IC}(c) = \frac{3}{12}$ $Cont_{\Psi}^{IC}(d) = \frac{7}{12}$. Operator Γ^{IC} returns $d < c < a = b$

Let us now consider a second measure which is more refined than Drastic measure. The measure is based on minimal inconsistency communication sets and the intuition is that the degree of inconsistency is proportional to the number of inconsistent subsets. The measure has to take care of our syntax-based perspective: to avoid syntactic biases, the measure considers the whole set of communications given by an agent.

Definition 6 (Minimal Inconsistent subsets measure). *Let Ψ be a set of communications on A and IC a consistent formula. The inconsistency measure I_{MI}^{IC} based on the number of minimal IC -contradictory subset is defined as:*

$$I_{MI}^{IC}(\Psi) = \left| \left(\bigcup_{a \in Ag(\Psi)} \langle a, \text{Report}(\Psi_a) \rangle \right) \perp IC \right|$$

Theorem 3. I_{MI}^{IC} is a syntax weak-independent inconsistency measure.

Then Γ_{MI}^{IC} is a good candidate to assess reliability. Let's illustrate this second measure on the same example.

Example 6. Consider agents a, b, c, d and propositional letters p, q, r . Consider $\Psi = \{ \langle a, p \wedge q \rangle, \langle b, q \rangle, \langle c, \neg q \rangle, \langle d, r \rangle \}$ and $IC = \neg r$. Then we get: $Cont_{\Psi}^{IC}(a) = \frac{1}{2}$ $Cont_{\Psi}^{IC}(b) = \frac{1}{2}$ $Cont_{\Psi}^{IC}(c) = 1$ $Cont_{\Psi}^{IC}(d) = 1$. Notice that the contributions of c and d are equal. According to these contributions Γ_{MI}^{IC} return either $d < c < a = b$ or $c < d < a = b$ depending on the deterministic choice.

6 Conclusion

This work proposes to assess the relative reliability of some information sources by analysing the consistency of information they report, whether it be the consistency of each single report, or the consistency of a report as regard to some trusted knowledge or the consistency of different reports together. We have given some postulates stating what the relative reliability preorder should be. Then we have introduced a generic operator for building such preorder which is parametrized by a function for measuring the inconsistency of the information reported. We prove that this generic operator agrees with the postulates.

This framework may be extended in several ways. First, inconsistency measures should deserve more attention. Recent work on this topic [7, 1] shows promising results such as giving a weight to the inconsistency itself. A second

issue concerns our key principle relying assessment considering only inconsistency: the more agent is connected to inconsistency, the less it is reliable. The reverse notion, might then also be considered for assessing the reliability. In other words, how can we “reward” an agent which is never inconsistent. The third issue concerns the one shot dimension of the assessment process: iteration should be possible and reliability assessment should then be viewed as a refinement process. In that case, the key issue is to set rationales for changing reliability from $a < b$ to $b < a$. Finally, it must be noticed that if one has already some partial information about the reliability of the agents (for instance, one knows that agent a is more reliable than b but has no idea about c reliability) then the process described in this paper is not applicable as is: in this case, we could achieve reliability assessment by combining that preorder with the one produced by the operator Γ^{IC} . For future work, we plan to study these aggregation operators.

As we can see the proposed framework offers numerous perspectives and our aim is to take advantage of its flexibility for going further.

Acknowledgements. We sincerely thank the anonymous reviewers whose questions helped us to improve the paper.

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