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# There are arbitrarily large minimal 2-pinning configurations 

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Let $\ell$ be a line and $\mathcal{C}$ a collection of disjoint convex sets in $\mathbb{R}^{d}$. We say that $\ell$ is a transversal to $\mathcal{C}$ if it intersects each of its members and that $\mathcal{C}$ pins $\ell$ if, in addition, no arbitrarily small perturbation of $\ell$ is a transversal to $\mathcal{C}$. We say that $\mathcal{C} k$-pins $\ell$ if for any $k$-flat $\Pi$ containing $\ell, \Pi \cap \mathcal{C}$ pins $\ell$ in $\Pi$ (so pinning and $d$-pinning are the same in $\mathbb{R}^{d}$ ). A minimal $k$-pinning configuration is a pair $(\ell, \mathcal{C})$ where $\mathcal{C} k$-pins $\ell$ but no proper subset of $\mathcal{C}$ does.

When $\mathcal{C}$ consists of disjoint balls, all its transversal that meet the balls in the same order make up a connected set in line space [1]. So, by continuity, $\mathcal{C}$ pins $\ell$ if and only if no other transversal to $\mathcal{C}$ realizes the same order as $\ell$. Since two lines span a space of dimension at most 3 , this implies that $\mathcal{C}$ pins $\ell$ if and only if it 3 -pins it. As a consequence, all minimal 3-pinning configurations of disjoint balls in $\mathbb{R}^{d}$ have size at most $2 d-1[1,2]$. Here, we prove that the situation is different for 2 -pinning:
Theorem 1. A minimal 2-pinning configuration of size $n \geq 6$ in $\mathbb{R}^{3}$ exists if and only if $n$ is even.
For simplicity we assume that the objects are smooth, e.g. balls. Our proof is essentially combinatorial:

1. Sign sequence. Let $\ell$ be an oriented line in $\mathbb{R}^{2}$ tangent to $O_{1}, \ldots, O_{n}$ in that order. Its sign sequence is a word in $\{-,+\}^{n}$ where the $i^{\text {th }}$ letter is + if and only if $O_{i}$ is on the right of $\ell$.
2. Encoding. We now return to 3 -space, assume that $\mathcal{C} 2$-pins $\ell$ and choose some arbitrary orientation on $\ell$. The tangent planes to the objects of $\mathcal{C}$ at their contact point with $\ell$ partition the space into $n$ sectors. For all planes containing $\ell$ lying in the same sector, the sign sequence of the associated planar configuration is the same (up to exchanging - and + ). We denote by $\sigma_{0}, \ldots, \sigma_{n-1}$ the sign sequences obtained successively as we go through all sectors around $\ell$.
3. Translating geometric properties. The family $\left(\sigma_{0}, \ldots, \sigma_{n-1}\right)$ corresponds to a geometric configuration $(\mathcal{C}, \ell)$ if and only if (1) there exists a permutation $\pi$ of $\{1, \ldots, n\}$ such that $\sigma_{i}$ differs from $\sigma_{i-1[n]}$ from the inversion of its $\pi(i)^{\text {th }}$ letter. Then, observe that $\mathcal{C} 2$-pins $\ell$ if and only if (2) every $\sigma_{i}$ contains an alternating triple. Last, notice that no proper subset of $\mathcal{C} 2$-pins $\ell$ if and only if (3) for any $1 \leq i \leq n$ there is some $t$ such that $\sigma_{t}$ loses all alternating triples if its $i^{\text {th }}$ letter is deleted.
4. Wrapping up. For odd $n$, rules (1)-(3) are incompatible. For even $n$, the sequence defined by

$$
\sigma_{0}=+-+^{n-2} \quad \text { and } \quad \pi=(4,1,6,3,8,5, \ldots, n, n-3,2, n-1)
$$

satisfies rules (1)-(3) and thus corresponds to a minimal 2 -pinning configuration.

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