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# New IPv6 Identification Paradigm: Spreading of Addresses Over Time

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**Abstract**—The identification of packet flows is a very important feature to provide security on the Internet. This flow identification is traditionally done by the well-known five tuple source IP address, destination IP address, transport layer protocol number and the two source/destination identifiers of transport layer protocols (named ports on UDP and TCP). Unfortunately, the IP source address is not reliable at all. However, we can use new security paradigms based on new IPv6 properties. In particular, IPv6 introduces a large address space. Our solution takes the benefit of this space with a high frequency rotation of IP addresses, that we call *spreading*. This spreading improves the security since only the sender and the receiver are able to generate and follow this temporal sequence. An attacker will not be able to successfully insert malicious packets into a flow or to initialize a flow. It protects against session initialization flooding and against attacks on established connections. In this paper, we describe the architecture of our solution and the protocol to initiate a connection and also performance evaluation of our spreading.

**Keywords**—IPv6; security; flow identification; spoofing.

## I. INTRODUCTION

### A. Flow identification in the current Internet

Flow identification is the base of some Internet mechanisms, like security (filtering of packets) and priority policies for Quality of Service. But since the Internet is a datagram network and IP is not a connection oriented protocol, the notion of flow is not explicit at the network layer. Each packet is independent, and two similar packets can follow two different routes. It is why a flow is defined in the RFC 2722 [1] as “an artificial logical equivalent to a call or connection”. It is “artificial”, and there are no easy ways to discriminate a flow in an IP network.

At the transport layer, the concept of flow is more natural. This is a mandatory function to sort packets, reassemble segments and detect errors, like the popular protocol Transmission Control Protocol (TCP) [2] does. To have this notion of flow, one can use Transport Layer addresses, named ports in case of TCP.

This is why we need a tuple of five elements in order to extract the notion of flow on the network. The first two are the source’s and destination’s IP addresses, directly available in the IP headers. The next one is the transport layer number, available in the field `next header` for a Internet Protocol version 6 (IPv6) packet without extensions. With the knowledge of the transport protocol, it is possible to parse the transport header

and also to read the port numbers. They complete the tuple, already filled with the three identifiers of IP header.

This well known five tuple is the basic identification of a flow. The identification can be more complex. For example, a stateful firewall will follow the TCP states of each connection, and it will discard packets if they do not follow the TCP standard.

### B. Address spreading benefits

These five members of the identification tuple ( $IP_{src}, IP_{dst}, NextHeader, Port_{src}, Port_{dst}$ ) are not authentic by nature. The source IP address and the source port especially can be manipulated easily by an attacker. If an attacker is able to send packets with a spoofed source IP address, he will be in good position to try TCP reset attacks, to inject packets on the destination network, to try a targeted attack to the destination, etc.

With IPv6, the large IP address space allows new security opportunities, like Cryptographically Generated Address (CGA) [3]. If all IPv4 solutions had to minimize the number of IP addresses in use, it is now possible to use a lot of addresses. Our solution provides security thanks to the *spreading* of addresses. In our solution, source and destination IP addresses of a flow are renewed frequently, according to a temporal sequence. If this sequence is known only by the sender and the receiver, it adds a new identification feature.

Since we only modify IP addresses, our solution is pretty simple. It does not need some complex encapsulation (like IPsec tunnel does), and can be followed by a firewall with the knowledge of a shared secret.

### C. Attacker model

Our attacker can inject some traffic with a spoofed source IP address. He can be on the transmission path and also able to read the legitimate traffic.

We do not try to protect against a rebuilding of a flow. This means that the attacker can use upper layers information like TCP ports and sequence numbers to rebuild the real flow. Our protection is against the spoofing: our spreader will recognize packets from the attacker. However, we provide a protection against correlation of flows, since we obfuscate addresses. An attacker can not guess the real source and destination addresses of a flow, and can not group several flows to one source or destination just with the help of the information available at the network layer.

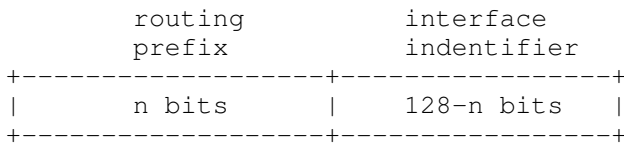


Figure 1: The two parts of IPv6 addresses

#### D. The two IPv6 address parts

An IPv6 address is divided in two parts, depicted in Figure 1. The routing part (usually called “prefix”) is not rewritable without connectivity issues, since such modification on the destination address will prevent upstream routers to send packets to the destination. In the same way, the source address has to be compliant with the anti-spoofing rules of the RFC 2827 [4] and could not be arbitrary rewritten.

There is no recommended size of prefix for a end network [5], but the length can not be more than 64 bits for compatibility with the IPv6 autoconfiguration. Indeed, the IPv6 Stateless Address Autoconfiguration (SLAAC) derives the second part of the address (named interface identifier) from the hardware MAC address to an identifier of 64 bits. A prefix larger than 64 bits breaks this autoconfiguration.

We choose to rewrite only the last 64 bits of each address, giving a total of 128 bits. Since 64 bits is the maximum size for a prefix compatible with autoconfiguration, this value will be compatible with almost all networks.

#### E. Related work

There is some previous work on dynamical IPv6 addresses. The nearest of our work is the “Moving target IPv6 Defense” publication [6]. This solution uses an User Datagram Protocol (UDP) tunnel to often rotate addresses and to encapsulate the real IP packet. An encryption is optional to protect the payload of the packet. The main drawback of this paper is the choice to make an encapsulation. It means that the solution has to fragment big packets to prevents Maximum Transmission Unit (MTU) problems and to reassemble it at the end of the tunnel. On the other hand, additional headers can have a bandwidth performance impact. This proposition does not use a temporal window for old address to prevent false positive.

An other idea is the publication “An Architecture for Network Layer Privacy” [7]. It uses an Site Multihoming by IPv6 Intermediation (SHIM6) extension to spread addresses over time. Since SHIM6 is an end-to-end solution, the end device has to allocate all possible addresses before using the connection. It is probably not desirable that a computer allocates 1000 addresses on one network interface, because it will send a lot of Neighbor Discovery packet.

An other paper [8] uses SHIM6 for a protection against deny of services. The idea is to switch from an address (under attack) to another one (not under attacks) for all connections already established.

#### F. Organization of this paper

In this paper, we first presents some problems of address spreading and our solution of architecture to overcome them

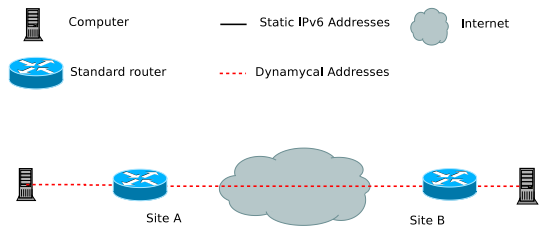


Figure 2: Spreading of addresses without extra device

in Section II. Second, we introduce principles of the spreading and some notations in Section III. We describe step by step the initialization of a connection in Section IV, that we complete with a description of packet processing on spreaders in Section V. We introduce the theoretical source of performance issues in Section VI. We close the paper with the description of performance evaluation in Section VII.

## II. LOCALIZATION OF THE SPREADER

### A. Difficulties to spread addresses on the end devices

The first idea of spreading is to generate and follow the address sequence on the end devices (see Figure 2). This strategy allows a end-to-end security, and the computer does not need to delegate the security to someone.

However, this solution implies an upgrade of the local router. The main issue arises if the end device does not get a delegated prefix, but share the local prefix with several end devices. The use of a lot of addresses will:

- flood the network with Neighbor Discovery packets. The router is not aware of the spreading, and can not know the link between the temporal sequence of IP addresses and the static MAC address;
- saturate the Neighbor table of the router. With frequently switching of addresses, the router will not be able to store all mappings between IP addresses and MAC addresses;
- introduce a latency at each IP address switching, due to the Neighbor Discovery.

To solve these issues, a solution is to patch the router to follow the sequence of IP destination addresses in the Neighbor table. Thus, the router does not need to know the sequence of source addresses received from the Internet, and is not able to insert a packet in a flow.

### B. The prefix delegation solution

With IPv6, we could have enough addresses to delegate one address prefix to each end device. It solves the problem of the mapping between MAC addresses and IP addresses, since the intermediate only send all IP packets matching a prefix to a MAC address.

In this case, the end device is in charge of Interface Identifiers management and can actually be seen as a “router”, and it is the same architecture at the network layer shown in Figure 3. The delegation of address prefixes is the best

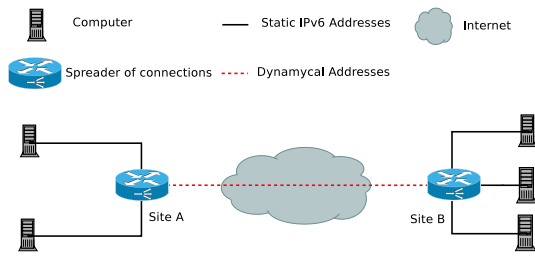


Figure 3: Architecture of the solution: spreader on the border of the local network

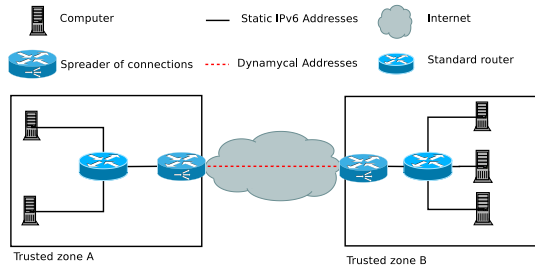


Figure 4: Architecture of the solution: spreader at the border of the trusted zone

architecture for the simplicity of the solution and from a security point of view. However, we need to provide a solution for a standard network without delegation of prefixes.

### C. Simplification with a spreader on the path

For a first approach, we propose to simplify the problem by adding some new “spreaders”, devices on the path of the communication. These devices are able to rewrite a flow of packets with stable addresses to a flow of packet with dynamical addresses. The spreader can be directly on the border of the network (see Figure 3) or at the border of the trusted zone (see Figure 4).

The first positive argument for this architecture is the simplicity to deploy the solution. An administrator does not need to upgrade and configure each end device, but can simply insert the spreader in the network. It has the same benefit than a modified router following the relation between IP and MAC address, and less drawback.

The second point is the ability of bad packets filtering. Since malicious packets consume resources, and can be send to simply saturate the bandwidth of the network, we have to discriminate malicious packets as soon as possible. With an introduction of the spreader at the border of the trust zone, we achieve efficiently this goal.

We choose this architecture for this paper to simplify concepts and experimentations.

## III. GENERAL PRINCIPLES

### A. Prerequisite of the solution

To enable the spreading, we need at least to configure two networks, for the mapping between the dynamical addresses

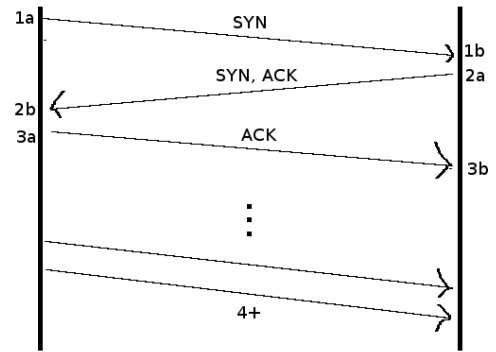


Figure 5: Digital timing diagram of a connection initialization

and stable identifiers.

This configuration is done by adding one spreader at the border of each network. The two spreaders have to share a secret, that an attacker can not guess. The communication of this secret is out-of-scope of this paper.

### B. Initialization of the spreader

The initialization of the spreader has to create a configuration for each compatible peer with a shared secret. This configuration contains the prefix list of the destination (to catch packets to be rewritten) and a function to derive cryptographic keys from the shared secret.

### C. Exchange of session data

One of our goals is to spread each flow of data with a unique sequence of addresses, make it more difficult for an attacker to group all flows of one end device. To provide it, both spreaders have to exchange session data at each flow initialization. There are several ways to accomplish it. The first one is to add several extras packets to initiate a context for each flow. It increases latency of connection initialization, and costs some bandwidth.

The second one is to add extra information on a real packet, like by adding one extra IPv6 extension header. Since this extension will be added by the spreader and not by the end devices, it could result in some maximal transport unit problem. Indeed, we can not add this header on a big packet, and have two choices. The first one is to fragment the packet on the spreader, IPv6 RFCs do not allow this solution. The second one is to send a “too big” error to the end devices, which reduce the performance of all packets for the session.

These solutions are not satisfying. We choose another one, our spreaders encode all information in IPv6 source and destination addresses, and do not add any extra data to packets payload. It limits the amount of exchangeable data, but does not have any cost of bandwidth or latency due to extra packets.

### D. Notations

The description of the protocol follows the same steps than a TCP connection initialization, depicted in Figure 5.

We introduce the following notation for the packet rewriting (summarized in Table I):  $P_A$  and  $P_B$  are the prefixes of networks for hosts A and B.  $IP_A$  and  $IP_B$  are the real IP addresses of host A and B, concatenation of prefixes and interface identifiers  $IID_A$  and  $IID_B$ .  $IP_{src}^n$  is the rewritten source IP address of the packet in step  $n$ , and  $IP_{dst}^n$  the destination IP address in the same step. Since we can not rewrite the prefixes,  $IP_{src}^n$  and  $IP_{dst}^n$  are concatenation of one stable prefix ( $P_A$  or  $P_B$ ) and one rewritten value.

TABLE I: IP ADDRESSES NOTATION

Steps	Local network		On the Internet	
	Real IP Source	Real IP Destination	Rewritten IP source	Rewritten IP Destination
$1a$ $\rightarrow 1b$ SYN	$IP_A$	$IP_B$	$IP_{src}^1 = P_A   IID_{src}^1$	$IP_{dst}^1 = P_B   IID_{dst}^1$
$2a$ $\rightarrow 2b$ SYN, ACK	$IP_B$	$IP_A$	$IP_{src}^2 = P_B   IID_{src}^2$	$IP_{dst}^2 = P_A   IID_{dst}^2$
$3a$ $\rightarrow 3b$ ACK	$IP_A$	$IP_B$	$IP_{src}^3 = P_A   IID_{src}^3$	$IP_{dst}^3 = P_B   IID_{dst}^3$

#### IV. STEP BY STEP INITIALIZATION

##### A. Initialization of a connection - first packet

1) *Symmetrical rewrite on spreaders*: The rewriting begins with step  $1a$ , when the spreader receives a packet with an destination IP address matching one of the prefix in the spreader configuration.

At the first packet of a connection, the local spreader computes new source and destination IP addresses with the help of a cryptographic function. We choose the Advanced Encryption Standard (AES) encryption, but it could be another one allowing encryption of blocks length of 128 bits.

The AES function takes as input a block of both interface identifiers of hosts A and B, and the actual key  $K(t)$  derived from the shared secret for the encryption. This AES function has an output of 128 bits, that the spreader divides in two blocks of 64 bits to replace last 64 bits of  $IP_A$  and  $IP_B$ .

$$IID_{src}^1 = AES(IID_A | IID_B, K(t))[0 - 63] \quad (1)$$

$$IID_{dst}^1 = AES(IID_A | IID_B, K(t))[64 - 127] \quad (2)$$

After the rewriting of addresses, the packet follows the standard routing and filtering process. This ends step  $1a$ .

On the destination spreader, stable addresses are recomputed in the same way (AES is a symmetric function) in the step  $1b$ . After the computing, the destination spreader checks the validity of the transport layer checksum (this checksum is mandatory for UDP and TCP with IPv6). If the checksum is valid, the packet follows the standard policy of routing and filtering.

If the checksum value is not valid, it can of course be a sign of transmission problem. One other possible cause of this invalid checksum is a try of an attacker to inject one packet in the network, with a spoofing of the source address. Indeed, IPv6 addresses are part of the checksum computation and if addresses after the second spreader are not the same than

addresses sent by the source device, it invalids the checksum. Since the attacker does not know the shared secret, he can not compute the AES encryption and the generated packet will be detected by the spreader.

In more details, the Table II depicts the status of the checksum around the Internet. It looks invalid between the two spreaders, but this is not a problem since nobody needs to have a look on the checksum in the path of the communication. On the contrary, we see in Table III that the checksum of a spoofed packet looks valid on the Internet, but invalid after the rewriting of the second spreader.

TABLE II: CHECKSUM VALIDITY FOR A REAL PACKET

Position	Checksum validity	Remarks
Local Network	Valid	Computed by the end device
Internet	Invalid	Addresses have been rewritten, it invalids the checksum
Remote Network	Valid	Addresses have been rewritten by the second spreader

TABLE III: CHECKSUM VALIDITY - SPOOFED PACKET

Position	Checksum validity	Remarks
Attacker's Network	Valid	Computed by the attacker
Internet	Valid	No rewriting if the attacker is not aware of protection
Remote Network	Invalid	Addresses have been rewritten by the second spreader

2) *Security analyze of the rewriting*: If the attacker is aware of this spoofing protection, he can try to guess the checksum modification added by the AES encryption of spreaders. The length of the checksum field is 16 bits, which it gives one chance on 65 536 to find the good one. This value is only valid for a short time and for a given address couple, the next value of  $K(t)$  at  $t+1$  will give another checksum modification implied by the AES encryption.

This security mechanism is not good enough to filter all packets of an attacker, and some packets can bypass this protection. But it is important to notice that if the checksum is valid, the attacker can not guess the rewritten addresses and can not know what is the rewritten destination address. The chance to successfully contact a real computer with a valid address is very low. Indeed, if we assume that the rewriting is fully random, an attacker has first to bypass the checksum (one chance on 65 536). If the checksum is valid, a targeted attack on a computer with an address on the remote network has one chance on  $2^{64}$  to reach the targeted computer, since an attacker can not guess the rewritten value after the AES decryption.

##### B. Initialization of the connection - reply of the remote

The goal of the addresses rewriting on the first packet is to protect an initialization of a connection by an attacker. For the

TABLE IV: REWRITING IN INITIALIZATION STEPS

Steps	Rewritten IID Source	Rewritten IID Destination
1a → 1b SYN	$IID_{src}^1 = AES(IID_A   IID_B, K(t))[0 - 63]$	$IID_{dst}^1 = AES(IID_A   IID_B, K(t))[64 - 127]$
2a → 2b SYN, ACK	$IID_{src}^2 = random()$	$IID_{dst}^2 = g(t, secret, IID_{src}^1)$
3a → 3b ACK	$IID_{src}^3 = g(t, secret, IID_{src}^1)$	$IID_{dst}^3 = g(t, secret, IID_{src}^2)$

next packets, we create a pseudo-random sequence for each flow of data, generated by a function  $g$ . This function  $g$  is a generator of a random temporal sequence, one example is a hash function like SHA1.

This creation begins in step 2a. To accomplish it, the second spreader rewrites the first reply packet of a client with a random value as IP source, and a value computed from the source IP address value in the first packet for the destination.

$$IID_{src}^2 = random() \quad (3)$$

$$IID_{dst}^2 = g(t, secret, IID_{src}^1) \quad (4)$$

$g$  takes as input the current time, the shared secret between networks and another value with the same size than an IP address. This rewriting introduces a random value for the sequence, but the flow is still easy to identify for both spreaders with the IP address destination sets to a value that the first spreader can recognize.

In step 2b, the first spreader recognizes the IP destination address  $IP_{dst}^2$  with the help of a context previously stored. This packet is an acknowledgment of the initialization, the spreading can now really begin. The spreader saves the value of the IP source (randomized in step 2a) and rewrites the source and destination IP addresses to the real stable value stored in the context.

It ends the second step. The first spreader is now sure of the initialization of the connection, and can use the random value to bootstrap a new random sequence.

### C. Initialization of the connection - Acknowledgment of the second spreader

The step 3a begins at the next packet sent by the device from the first network. Both source address and destination address are now spread with:

$$IID_{src}^3 = g(t, secret, IID_{src}^1) \quad (5)$$

$$IID_{dst}^3 = g(t, secret, IID_{src}^2) \quad (6)$$

In step 3b, the second spreader recognizes the couple with the help of the stored context. This packet is an acknowledgment of the random value send in step 2a, and the second spreader is now aware of the success of the initialization. Steps of initialization are summarized in Table IV.

### D. Rewriting during the life of the connection

After the step 3b, both spreaders follow the same sequence of rewriting with  $g(t, secret, IID_{src}^1)$  and  $g(t, secret, IID_{src}^2)$ . The rewriting is symmetrical and both end devices receive stable addresses. An attacker can not inject some traffic since he does not have the knowledge of the next addresses in use.

### E. Definition of a flow, division of connections in several temporal sequences

Our goal is to create one sequence of address for each flow of data. By flow, we mean a sequence of packets where informations of upper layers are enough for an attacker to correlate and to rebuild the sequence.

The first trivial idea is to make one flow for each couple of  $IP_{src}, IP_{dst}$ . It does not take a lot of resources, but do not prevent correlation if more than one flow are send between devices. Nevertheless, it can be desirable to obfuscate this information.

IPv6 introduces a new header field to give information about a flow of packets: the flow label field [9]. It is a 20 bits header, that can be used for Quality of Service or other usages (the RFC 6294 [10] tries to give of survey of usages). This flow label can be rewritten on the path of the communication, and is not part of the checksum. Other usages than Quality of Service are allowed, which means that we respect standard specifications [9] with our proposition.

This is why we define a flow by a tuple of  $(IP_{src}, IP_{dst}, flowlabel)$ . If the end device set a different value for two flows, it will be spread into two different sequences. The flow label is set by the end device itself, which has enough information to know if a sequence of packets should not be grouped with another connection.

## V. DETAILED PROCESSING ON SPREADERS

Our description of the protocol in Section IV is the ideal situation. We do not have any loss of packets, and the end devices use the TCP protocol. It helps to understand how our protocol works, since it follows the same handshake mechanism than TCP.

But we have to be robust against loss of packets and retransmission of data. In the same way, we have to be compatible with transport protocols, like UDP, where a connection does not follow a rigorous initialization procedure, or ICMP with short session like a simple Echo request. We detail in this section the processing of packets in spreaders, which actually support loss of packets and is transport layer protocol independent.

### A. Detailed steps of packets processing (outgoing packets)

The processing of packets from the local network to the Internet is depicted in Figure 6. For each outgoing packet, the spreader extracts the tuple  $(IP_{src}, IP_{dst}, flowlabel)$  (step 1). It checks if one context already exists (step 2) and if we already received an acknowledgment (step 3). If both conditions are true, we are in the case of an established connection and we can rewrite  $IP_{src}$  and  $IP_{dst}$  with the help of the function  $g$

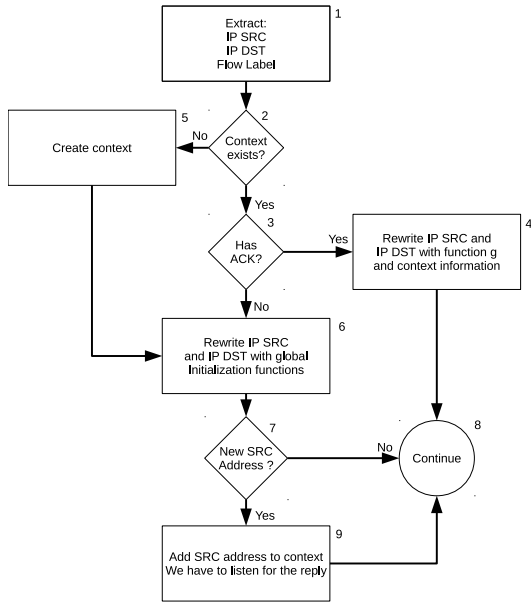


Figure 6: Processing of outgoing packet (to the Internet).

and information stored in the context. After that, the packet continues the standard processing.

If the context does not exist, we have to create one. It contains the real IP source address and real IP destination address, as well as the flow label value (step 5). We move to step 6, the case of the context exists but we do not have yet any acknowledgment. We have to use AES encryption with  $K(t)$  to rewrite IP addresses.

Since the rewritten  $IP_{src}$  is used as parameter for a reply from the remote network, we have to store it (step 7 and step 9). We can store more than one address if we send several packets before we receive any acknowledgment. The packet follows afterwards the standard packet processing.

### B. Detailed steps of packets processing (incoming packets)

The Figure 7 depicts the packet processing for all packets from the Internet. The processing of incoming packets begins with the extraction of the couple  $(IP_{src}, IP_{dst})$ . We do not extract the flow label, this value can not be trusted outside of the local network. This flow label will be rewritten to an internal value to make the future flow identification of local packets going to the Internet. The goal is to know if a context already exists for this connection (step 2 and 3). If this is the first packet for this context, it is an acknowledgment of the initialization and we have to change the status of the context (step 8). We end the processing with the rewriting of dynamical addresses to stable addresses and we return the packet for standard processing (step 6).

If we do not have a context, we have to try a decryption of IP addresses with the AES function and  $K(t)$  in step 7. This decryption is followed by the computation and the verification of the transport layer protocol checksum in step 9. A bad checksum implies to drop the packet, since it is probably a try of an attacker to send a packet with a spoofed address. If

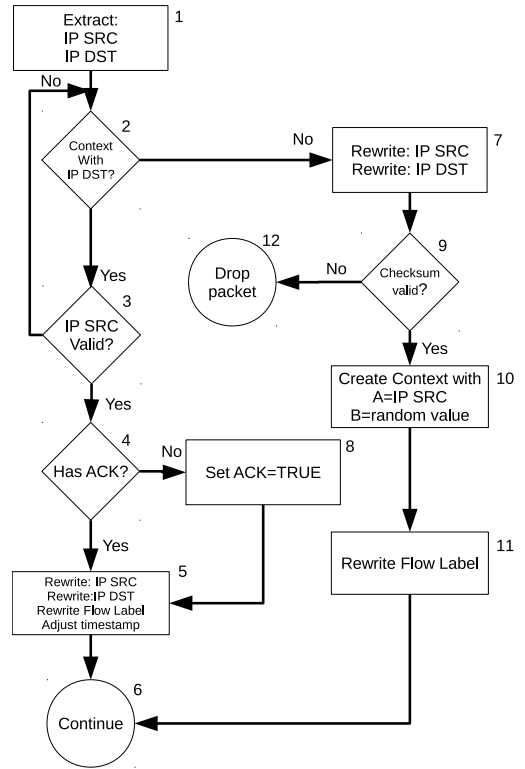


Figure 7: Processing of incoming packet (from the Internet).

it is valid, we initiate a context (step 10) and we return the rewritten packet for standard processing.

## VI. THEORETICAL LOSS OF PACKETS

Our spreading solution drops packets if they are not following the sequence of addresses. This spreading protects against an attacker, but valid packets sent by the real device can be dropped. Indeed, the latency in the network is the main problem: if a packet is too long in transmission across the network, it will be drop by the receiver, since he had already switched to the next addresses.

The second source of problems is the time desynchronization between two spreaders: if the clocks are not synchronized, a valid packet will be detected as spoofed even if the sender is not an attacker.

In this section, we describe the theory of this packet loss. We first estimate the loss of packets in case of latency in the network with perfectly synchronized spreaders. Second, we add the problem of time desynchronization between spreaders. Next, we explore the consequences for a simple ICMP echo request/echo reply communication. We conclude with solutions to mitigate both problems.

### A. Latency effect

The latency is the time needed for a packet to go from the source to the destination. The latency can be less than one millisecond on a local network (LAN), and several seconds between both points on the Internet. If we assume that the latency is stable for all packets and is the same in both

directions, it is easy to estimate the proportion of the packet loss. All packets sent at the end of the lifetime of one address will be drop by the receiver spreader. The duration of this black hole in the communication is exactly the value of the latency. If we consider that packets are uniformly send over time, we have them a packet proportion loss of:

$$loss = \frac{latency}{lifetime} \quad (7)$$

For example, with a configuration of 1 second for the lifetime of addresses, with 100 ms of latency on the network, 10% of packets will be dropped in both direction of the transmission between the spreader A and the spreader B.

### B. Desynchronization effect

It is not easy to perfectly synchronize two computers on the Internet. Even with the Network Time Protocol, clocks of computers are not perfect and there are always a little desynchronization. In one way of the transmission, this desynchronization is good, since it reduces the observed latency between both computers. In the other direction, the latency is added to the latency and it implies a longer duration of black hole for the communication.

If we assume that the spreader A is desynchronized in the future with spreader B, the loss in the direction A to B is:

$$loss_{AB} = \frac{latency + desync}{lifetime} \quad (8)$$

In the other direction from B to A, the loss has been reduce to:

$$loss_{BA} = \frac{|latency - desync|}{lifetime} \quad (9)$$

The perfect case for this direction is when the latency is equal to the desynchronization, there are no more packet loss in this direction. If the desynchronization is bigger than the latency, some packets come too early to the spreader A (A did not yet switch to the “new” address) and packets are dropped.

With the same configuration of 1 second for the lifetime of addresses, with 100 ms of latency on the network, and a desynchronization of 50 ms, 15% of packets are dropped between A and B (150ms of black hole) and 5% between B and A.

### C. ICMP echo request/echo reply communication

The evaluation of packet loss in both directions is not enough to evaluate the impact on a communication. On the Internet, a unidirectional payload transmission is very unusual. The most popular protocol TCP sends a lot of acknowledgment packets even for a unidirectional transmission, and a simple ICMP echo request is replied with an echo reply packet.

Our loss of packets is not distribute like a standard network error, each packet does not have a probability of failure, but the connection seems to be broken for a short duration, in one or two ways.

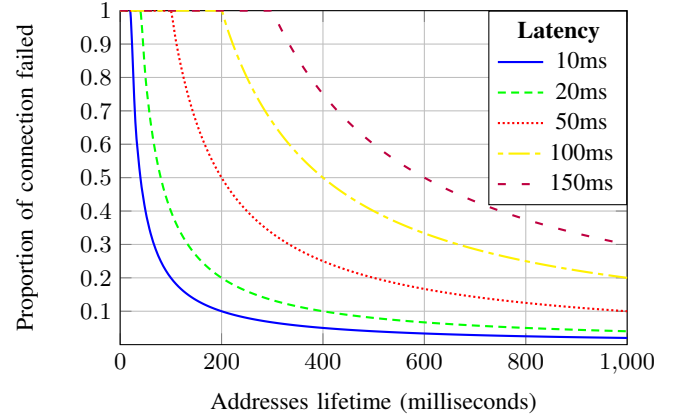


Figure 8: ICMP echo failure rate in function of latency

For example, in case of ICMP echo transmission between spreaders A and B, the total duration of one of packet echo reply/request will dropped is

$$latency * 2 + desync \quad (10)$$

We plot in Figure 8 the theoretical ICMP echo request/reply transmission failure for a given latency in function of address lifetime, without any desynchronization of spreaders.

### D. Temporal windows one old and next addresses

To reduce the loss of packets, it is possible to accept the old address of time  $t-1$  in a temporal windows where both current address  $t$  and  $t-1$  addresses are accepted. With a temporal windows larger than the latency, no packets are dropped by synchronized spreaders.

In case of desynchronization smaller than the latency, we have to add this desynchronization duration to the temporal windows to accept all packets send in the communication. A spreader can not know if a packet is delayed due to the latency or due to a desynchronization problem.

If the desynchronization is larger than the latency, a spreader will receive packets to early. To solve it, we can add in the same way a temporal windows where both addresses of  $t$  and  $t+1$  are valid. If a spreader receives a lot of packets on the  $t+1$  address, it is a sign of desynchronization and it could help to resynchronize both spreaders. We present our experimental results on temporal windows to accept old and future addresses in Section VII-D.

## VII. PERFORMANCE TESTS

### A. Test beds

1) *Implementation:* Our implementation is based on a Netfilter module for Linux. It follows steps explained in Section V. The implementation supports configuration of the validity lifetime of one address as well than temporal windows for packets out of the current time sequence.



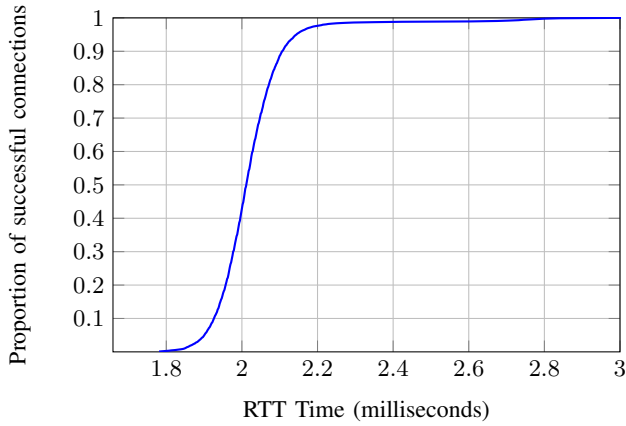


Figure 9: RTT of UDP echo requests on the LAN

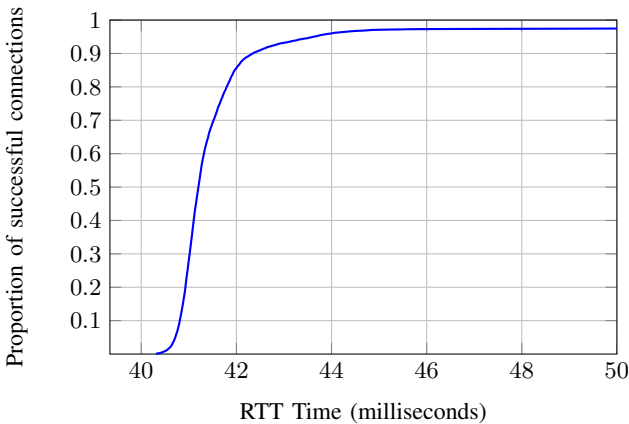


Figure 10: RTT of UDP echo requests on the 6to4 network

2) *Networks tests:* We have done several experimentations on several test beds, and this paper presents results from two of them. Our first test bed is the ideal one, in a LAN. The typical round-trip delay time (RTT) is around 2 milliseconds (ms), there are no loss or desequencing of packets. The Figure 9 depicts the cumulative distribution function of the RTT.

The second one is between a server in Germany using a 6to4 tunnel and a server with native IPv6 connectivity in France. The network has a poor quality: there are some natural packet loss (around 1%) and desequencing (around 0.5% on high network load). In our preliminary tests, the 6to4 tunnel uses to be less congested in the night, and we ran our tests in the night to avoid some random congestion issue. The Figure 10 depicts the cumulative distribution function of the RTT.

All devices of our networks are time synchronized on the same Network Time Protocol server. It does not provide a perfect synchronization.

## B. Spreading consequences

1) *Tests description:* To evaluate the consequences of our spreading, we ran for each network three tests. The first one

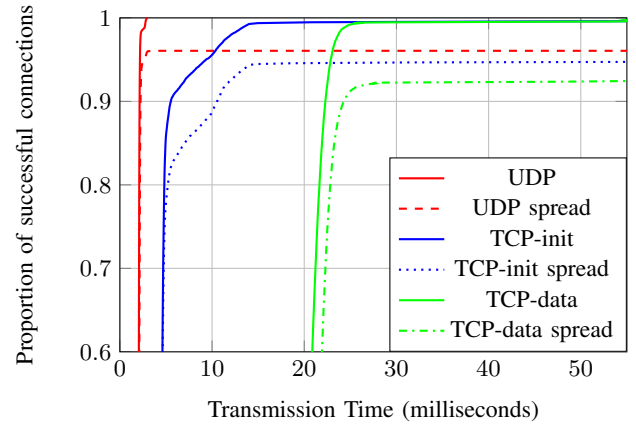


Figure 11: Consequences of a simple spreading on the LAN

send a standard UDP echo packet, it provides a good evaluation of the network quality for the loss of packets and the latency. The second one is a simple TCP handshake initialization, without data transfer. The last one is a TCP connection with data transfer (65535 bytes). Since TCP is the most popular protocol on the Internet and that the test involves transfer of data, it is the best test to evaluate the user experience on a network with spreading. We set a timeout of 4 seconds on both TCP tests, and we consider it loss after this time.

2) *Simple spreading:* Since computers are not perfectly synchronized and due to the network latency, our spreading implies some loss of packets and it has consequences on the network quality. We set first a lifetime of one second for each address, without any temporal windows for old and future addresses. The Figure 11 depicts the results on the LAN and Figure 12 on the 6to4 network.

On the LAN, there are no loss of UDP packet without spreading. With spreading enabled, the percentage of failure is around 4%. UDP does not provide retransmission of data and the proportion of success does not increase with the time.

The TCP handshake needs three packets to be completed, and the opening time of the majority of TCP connection is consistent with the UDP test. Some openings are delayed, and will be successfully completed with retransmission. They are not displayed on Figure 11, but we observe the same kind of results that are shown in the 6to4 network.

On the 6to4 network, loss of packets has big effect on TCP performance. For the opening of the connection, we see some steps corresponding to standard time of Linux retransmission strategy for TCP.

3) *High frequency addresses switching, consequences on loss of packets:* Of course, the lifetime of addresses has a big impact on the quality of the connection. We tried lifetime values between 50 ms and 2 seconds, and we summarize the percentage of packet loss in Table V.

We compare it to the theoretical result of Section VI-C in Figure 13. For the 6to4 network, the typical latency is around 21 ms, and we measure a desynchronization of 8 ms at the end of the experiment, it gives us around 50 ms of black hole

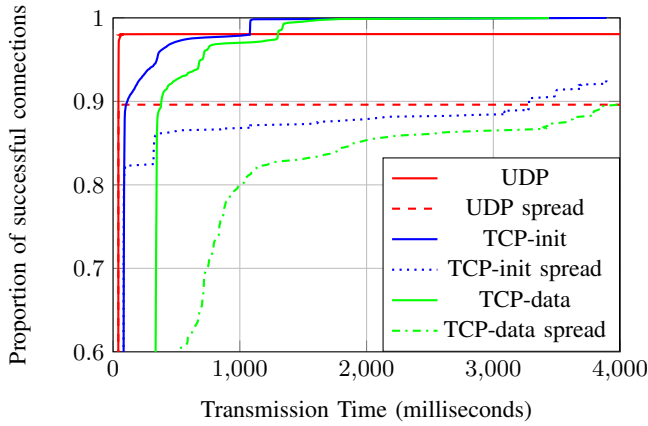


Figure 12: Consequences of a simple spreading on 6to4

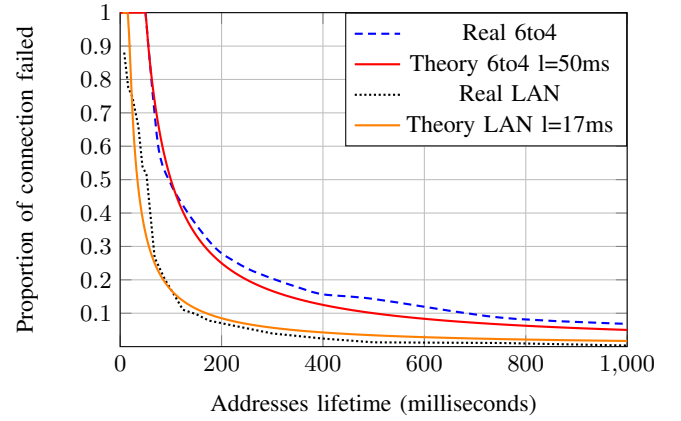


Figure 13: Proportion of failed ICMP echo transmission

TABLE V: PROPORTION OF PACKET LOSS

Lifetime	LAN	6to4
50 ms	51%	100%
75 ms	13.2%	60%
100 ms	10%	48.8%
150 ms	9.7%	36.5%
200 ms	5.7%	28%
300 ms	3.9%	20.4%
500 ms	1.3%	14.3%
750 ms	1.1%	8.7%
1000 ms	0.3%	6.8%
2000 ms	0.1%	5.1%

for the transmission. Our experimental result is very close to this theoretical result. But since the 6to4 network has some natural packet loss, there are more failure than expected when the spreading effect decrease.

On the LAN, the latency is small, and the desynchronization is the main issue. We measure a desynchronization between 0 and 30 ms, it is not stable in time of experiences. We took the middle value to plot the theoretical value with a black hole of 17 ms at each address switching.

#### C. Delayed packets: temporal window for the last old address

To prevent loss of delayed packets, we add a temporal window for the old address. In this temporal windows, the old and the current addresses of a sequence are accepted by the spreader. It is very efficient to decrease the loss of packets on the 6to4 network, like we can see in Figure 14.

With a temporal windows larger than the sum of the latency and the desynchronization of the network, we get the same performance than without spreading.

As depicted in Figure 15, it is not enough to prevent loss of packets on the LAN. The desynchronization is bigger than the

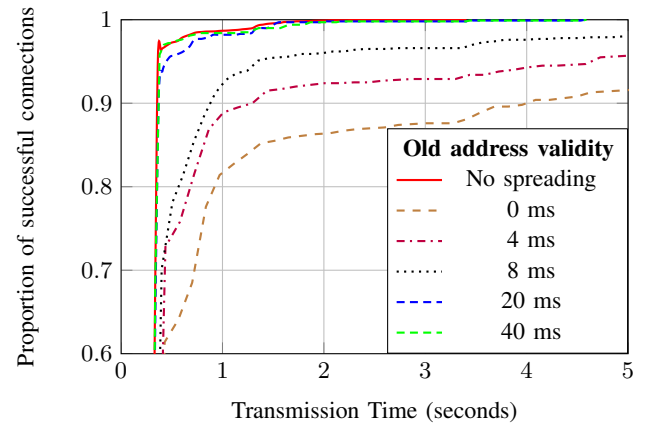


Figure 14: TCP test with transmission of data on the 6to4 network in function of temporal windows for the old address

latency and some packets come to early for a spreader. In this case, we need to accept the previous address on the spreader desynchronized in the future as well than the next address on the other spreader.

#### D. Desynchronization: temporal window for next address

To solve the desynchronization issue on the LAN, we add a temporal windows for the next address. During this temporal windows, both actual address and next address of the sequence are accepted by the spreader.

We plotted the results of UDP test in Figure 16, with a lifetime of 200 ms and a temporal windows for the old address of 60 ms.

Adding a temporal windows for the next address is not enough to avoid any loose of packet. We need to accept the old address on the spreader with a clock ahead of the real time. We wrote in Table VI the loss of packets with respect of both temporal windows for the old and the future address on the LAN. Since the desynchronization is not stable in the time of the experience, some values can be confusing. The loss can decrease if we increase the temporal windows. However, if

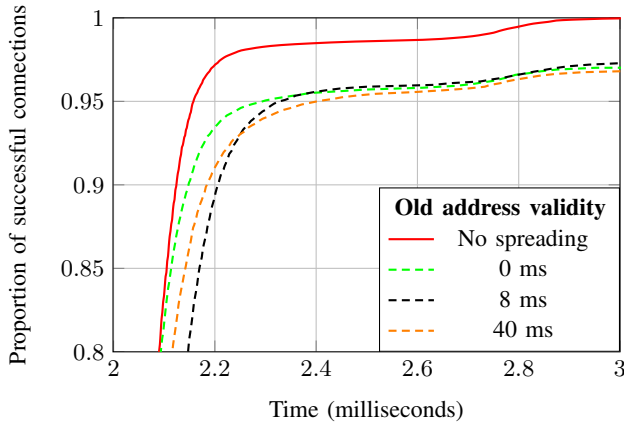


Figure 15: UDP test in function of temporal windows for the old address on the LAN

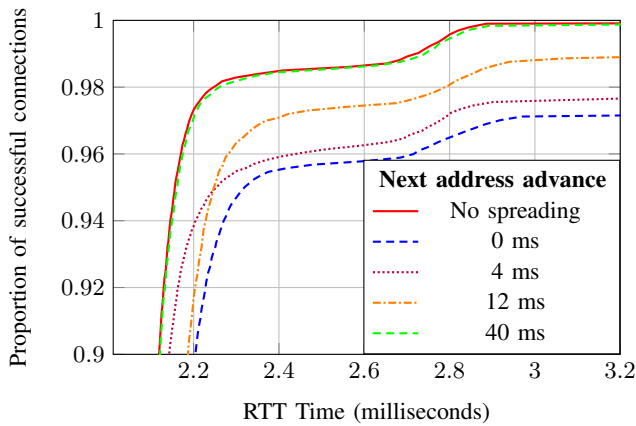


Figure 16: Next address temporal windows on the LAN

both temporal windows are larger than the desynchronization, there are no more loss of packets. A value of 64 milliseconds is enough here.

## VIII. CONCLUSION

The spreading of addresses is an innovative and new solution to identify a connection. This is a new mechanism to protect against spoofing attack. Our spreading protects against

TABLE VI: INFLUENCE OF NEXT AND OLD ADDRESSES ON PACKET LOSS

% loss		Next address temporal windows (ms)							
		0	4	8	12	20	40	64	80
Old address validity	0	3.5	5	3.9	3.3	3.3	4.8	3	4.9
	4	6	3.3	0.8	0.9	0.5	1.3	1.1	1.5
	8	2.7	2.3	0.6	0.7	0.7	0.1	0	0
	12	6.4	3	2	0	0	0	0	0
	20	3	2.3	0.5	0	0	0	0	0

initialization of a connection from an attacker, as well than injection of packet inside a established connection.

We described a complete protocol to securely initiate connection between spreaders, with one initialization of temporal sequences of addresses per flow. We did a step by step description of the spreader internal functionality, and we explain the theoretical loss of packets without any temporal windows.

With the use of temporal windows for the old address, we can protect against false positive detection of packets due to the network latency. With the use of a temporal window for the next address, we protect our solution against desynchronization of devices. We can use this information to resynchronize spreaders without external source of time.

Thanks to these temporal windows, we can achieve a very high frequency of address switching. An address is valid only for several duration of the latency in the network.

In the future works, we will evaluate algorithms to generate sequence of addresses. To simplify the work of the network administrator, the auto-configuration and the exchange of the secret between the spreaders have to be considered.

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