

Quantifiers for randomness of chaotic pseudo-random number generators

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We deal with randomness quantifiers and concentrate on their ability to discern the hallmark of chaos in time series used in connection with pseudo-random number generators (PRNGs). Workers in the field are motivated to use chaotic maps for generating PRNGs because of the simplicity of their implementation. Although there exist very efficient general-purpose benchmarks for testing PRNGs, we feel that the analysis provided here sheds additional didactic light on the importance of the main statistical characteristics of a chaotic map, namely (i) its invariant measure and (ii) the mixing constant. This is of help in answering two questions that arise in applications: (i) *which is the best PRNG among the available ones?* and (ii) *if a given PRNG turns out not to be good enough and a randomization procedure must still be applied to it, which is the best applicable randomization procedure?* Our answer provides a comparative analysis of several quantifiers advanced in the extant literature.

Keywords: random number; statistical complexity; recurrence plots; excess entropy; rate entropy; permutation entropy

1. Introduction

Chaos theory started more than 30 years ago and changed our world view regarding the role of randomness and determinism. As the statistical characteristics of chaotic systems were better understood (Lasota & Mackey 1994; Brown & Chua 1996; Beck & Schlögl 1997; Setti *et al.* 2002), a wide variety of situations emerged in which chaos, instead of stochastic systems, became a ‘controller of noise’.

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Chaos illustrates the rather striking fact that complex behaviour arises from simple rules when nonlinearities are present. Since simple chaotic maps are capable of generating stochastic-like signals, implementations based on chaotic systems are usually less involved than those based on more complex algorithms (Stojanovski & Kocarev 2001; Kocarev & Jakimoski 2003). One tries to apply this notion to generate pseudo-random number generators (PRNGs) because random numbers are widely used not only in cryptography and Monte Carlo applications but also in less obvious applications (Pecora *et al.* 1990; L'Ecuyer 1994; Kocarev & Parlitz 1995; Hidalgo *et al.* 2001; Fernandez *et al.* 2003). We mention just a couple of them. (i) In spread spectrum techniques, a binary signal is mixed with a random number sequence, to spread the spectrum over a wider frequency range. Using different random number sequences, it is possible to share a communication channel among several users (Mazzini *et al.* 1997; Dinan & Jabbari 1998; Shan *et al.* 2006; De Micco *et al.* 2007). Reduction of electromagnetic interference is another important benefit of the spread spectrum effect (Setti *et al.* 2000; Callegari *et al.* 2002). (ii) Consider a low-frequency signal immersed in a high-frequency digital noise. Sampling at time intervals defined by a random number sequence, the resultant signal becomes filtered without using any coil or capacitors that are expensive, especially in power systems (Petrocelli *et al.* 2007).

Truly random numbers are not attainable from computers, and it is unlikely that we will ever be able to get them from 'natural' sources, since one commonly assumes that any system is governed by underlying physical rules and consequently it is deterministic. A successful strategy to build up a PRNG is to start with the time series of a simple nonlinear chaotic map and to apply to it an adequate *randomizing procedure* so as to 'heighten/boost' its stochastic nature. Such strategy requires a quantitative evaluation of the improvement achieved after effecting the procedure. González *et al.* (2005) used the statistical complexity measure originally proposed by López-Ruiz *et al.* (1995) and later modified by Lamberti *et al.* (2004) to quantify the effectiveness of such randomizing modus operandi when applied to a Lorenzian chaotic system. It was also shown there that a widely employed course of action—the mixing of two chaotic signals—is not effective in this respect, contrary to what one might expect. In this vein, it is important to note that general-purpose tests available in the open literature (<http://csrc.nist.gov/rng/>; <http://stat.fsu.edu/pub/diehard/>; <http://www.iro.umontreal.ca/~simardr/random.html>) are not designed taking into account the particular characteristics of a chaotic map. Instead, one can appreciate in Rosso *et al.* (2007a) the fact that the deterministic nature of chaotic dynamics leaves special manifestations in the associated time series, which can be revealed only by recourse to adequate statistical quantifiers.

De Micco *et al.* (2008) randomized chaotic maps by means of two different randomizing procedures, namely *discretization* and *skipping*. The idea of concocting an 'information' plane, called the entropy-complexity plane, was advanced in order to use it as a means to ascertain the effectiveness of each of these two modi operandi. The main difference of this information plane from other complexity-entropy diagrams is the joint use of two different probability distribution functions (PDFs) in it, both associated to the pertinent time series.

Other important tools at our disposal are to be mentioned as well. In a recent and excellent report, Marwan *et al.* (2007) reviewed applications of so-called recurrence plots for a wide variety of fields and endeavours. They also proposed

several measures to quantify the recurrence plots' characteristics (Marwan *et al.* 2007). Additionally, two useful information-theoretic quantifiers of randomness, the *rate entropy* and the *excess entropy*, were proposed in Feldman *et al.* (2008) as coordinates of a complexity-entropy diagram.

In the present work, we explore combinations of all the above-mentioned quantifiers with the purpose of answering the following questions: (i) *among several chaotic maps, just which is the one that generates the best time series?* and (ii) *which is the best strategy—discretization or skipping—to randomize a given chaotic time series?* The ensuing quantifiers testing will be performed by means of two representative chaotic maps (and their iterates).

The paper is organized as follows: the statistical properties of a chaotic map and its iterates are reviewed in §2. Section 3 describes each of the analysed quantifiers. Section 4 deals with results for two representative maps and, finally, conclusions are presented in §5.

2. Statistical properties of a chaotic map

Let f be a chaotic map on the interval $[0, 1]$. Suppose the map has an invariant measure $\mu(x)$. Then the map is *ergodic* if for any integrable test function $Q(x)$, and for an arbitrary initial condition x_0 (up to a set of zero μ -measure), the time average is equal to the ensemble average

$$\bar{Q} = \langle Q \rangle. \quad (2.1)$$

Equation (2.1) is a consequence of the famous Birkhoff ergodic theorem (Cornfeld *et al.* 1982). *Mixing* is an even stronger requirement than ergodicity. A map is called '*mixing*' if any smooth initial probability density $\rho(x)$ converges to the invariant measure $\mu(x)$ after enough successive iterations. Mixing implies ergodicity. The reverse, however, is not true (Beck 1990).

There exists an equivalent definition of mixing via *correlation functions*. Let $\phi_1(x)$ and $\phi_2(x)$ be two integrable test functions and define the generalized correlation function of the map f by

$$C(\phi_1, \phi_2, n) = \lim_{J \rightarrow \infty} \frac{1}{J} \sum_{j=0}^{J-1} \phi_1(x_{j+n}) \phi_2(x_j) - \langle \phi_1 \rangle \langle \phi_2 \rangle, \quad (2.2)$$

where

$$\langle \phi_i \rangle = \lim_{J \rightarrow \infty} \frac{1}{J} \sum_{j=0}^{J-1} \phi_i(x_j). \quad (2.3)$$

The map is mixing if, for arbitrary ϕ_1 and ϕ_2 ,

$$\lim_{n \rightarrow \infty} C(\phi_1, \phi_2, n) = 0. \quad (2.4)$$

Let us stress that it is not easy to prove that f is a *mixing map* because the mixing condition given in equation (2.4) must be fulfilled for *arbitrary* test functions. Formally, every mixing map f has an associated

Perron–Frobenius operator \mathcal{L} (Beck 1990) that determines the time evolution of any initial density $\rho_0(x)$ towards the invariant measure $\mu(x)$

$$\rho_{n+1} = \mathcal{L}[\rho_n]. \quad (2.5)$$

The explicit formal expression for the Perron–Frobenius operator for a one-dimensional map f is given by (Beck 1990)

$$\mathcal{L}[\rho_y] = \sum_{x \in f^{-1}(y)} \frac{[\rho_0(x)]}{|f'(x)|}. \quad (2.6)$$

This operator \mathcal{L} has a set of eigenfunctions $\psi_\alpha(x)$ and eigenvalues η_α . The invariant measure $\mu(x)$ is the eigenfunction corresponding to the largest eigenvalue $\eta_0 = 1$. The full set of eigenfunctions and eigenvalues may be used as a basis to express any density

$$\begin{aligned} \rho_0(x) &= \sum_{\alpha} c_{\alpha} \psi_{\alpha}(x), \\ \rho_n(x) &= \mathcal{L}^n \rho_0(x) = \sum_{\alpha} \eta_{\alpha}^n c_{\alpha} = c_0 \psi_0(x) + R_n. \end{aligned} \quad (2.7)$$

The eigenvalue with the second largest absolute value, η_1 , has a ‘distinguished’ physical meaning: it is related with the *mixing constant* r_{mix} that governs the relaxation of *exponentially mixing* maps

$$|R_n| \sim |\eta_1|^n \sim \exp\left(-\frac{n}{r_{\text{mix}}}\right). \quad (2.8)$$

The invariant measure $\mu(x)$ gives the histogram of the time series, and the ideal PRNG must have $\mu(x) = \text{const}$. The mixing constant r_{mix} gives the transient characteristic time (Mazzini *et al.* 1997; Rovatti *et al.* 2004*a,b*; De Micco *et al.* 2007; Petrocelli *et al.* 2007), and its ideal value is $r_{\text{mix}} = 0$. In many applications of PRNGs, both the invariant measure and the mixing constant are relevant.

The analytical expression of the invariant measure $\mu(x)$ of a given map f is usually not known. Exceptions are the logistic map in full chaos and the piecewise-linear maps. The mixing constant r_{mix} has been analytically obtained only for piecewise-linear maps. For other maps, it must be numerically obtained by means of a piecewise-linear approximation of the map (Lasota & Mackey 1994; Beck & Schlögl 1997).

It is then obviously convenient to have quantifiers for measuring the uniformity of the invariant measure $\mu(x)$, and the mixing constant r_{mix} , for any chaotic map. These quantifiers are useful to compare time series coming from different chaotic maps and also to assess the improvements produced by *randomization procedures*.

It is possible to show that the invariant measure of f^d is identical to the invariant measure of f . Also, the mixing constant r_{mix} for f^d is lower than the mixing constant for f . The iteration of a map is one of the randomization procedures proposed in the literature, being used to diminish r_{mix} . This procedure is also known as *skipping* because iterating is tantamount to skipping values in the original time series, which does not change $\mu(x)$ and, consequently, is not conducive to a randomization of chaotic maps with $\mu(x) \neq \text{const}$. In this paper, we will use ‘skipping’ as a method of quantifier analysis.

3. Quantifiers for the invariant measure and mixing constant

In this section, we review several quantifiers proposed for measuring the main statistical properties of chaotic PRNGs. The quantifiers are classified according to their origin into three classes: (i) quantifiers based on information theory (López-Ruiz *et al.* 1995; Lamberti *et al.* 2004; Rosso *et al.* 2007a), (ii) quantifiers based on recurrence plots (Eckmann *et al.* 1987; Marwan *et al.* 2007), and (iii) quantifiers based on intrinsic computation (Feldman *et al.* 2008).

(a) Quantifiers based on information theory

They are appropriate functionals of the PDF. Let $\{x_i\}$ be the time series under analysis, with length M . There are infinite possibilities to assign a PDF to a given time series, a subject that will be given due consideration below. In the meantime, suppose that the PDF is discrete and is given by $P = \{p_i; i = \dots, N\}$. One then defines various quantities, as follows.

(i) *Normalized Shannon entropy* $H[P]$. Let $S[P]$ be the Shannon entropy

$$S[P] = - \sum_{i=1}^N p_i \ln(p_i). \quad (3.1)$$

It is well known that the maximum $S_{\max} = \ln(N)$ is obtained for $P_e = \{1/N, \dots, 1/N\}$, that is, the uniform PDF. A ‘normalized’ entropy $H[P]$ can also be defined in the fashion

$$H[P] = S[P]/S_{\max}. \quad (3.2)$$

(ii) *Statistical complexity measure*. A full discussion about statistical complexity measures exceeds the scope of this presentation. For a comparison among different complexity measures, see the excellent paper by Wackerbauer *et al.* (1994). In this paper, we adopt the definition of the seminal paper of López-Ruiz *et al.* (1995) with the modifications advanced in Lamberti *et al.* (2004) so as to ensure that the concomitant SCM version becomes (i) able to grasp essential details of the dynamics, (ii) an intensive quantity, and (iii) capable of discerning both among different degrees of periodicity and chaos (Rosso *et al.* 2007a). The ensuing measure, to be referred to as the intensive statistical complexity, is a functional $C[P]$ that reads

$$C[P] = Q_J[P, P_e] \cdot H[P], \quad (3.3)$$

where Q_J is the ‘disequilibrium’, defined in terms of the so-called extensive Jensen–Shannon divergence (which induces a squared metric; Lamberti *et al.* 2004). One has

$$Q_J[P, P_e] = Q_0 \cdot \{S[(P + P_e)/2] - S[P]/2 - S[P_e]/2\}, \quad (3.4)$$

with Q_0 a normalization constant ($0 \leq Q_J \leq 1$) that reads

$$Q_0 = -2 \left\{ \left(\frac{N+1}{N} \right) \ln(N+1) - 2 \ln(2N) + \ln N \right\}^{-1}. \quad (3.5)$$

We see that the disequilibrium Q_J is an intensive quantity that reflects on the system's 'architecture', being different from zero only if there exist 'privileged', or 'more likely' states among the accessible ones. $C[P]$ quantifies the presence of correlational structures as well (Martin *et al.* 2003; Lamberti *et al.* 2004). The opposite extremes of perfect order and maximal randomness possess no structure to speak of and, as a consequence, $C[P]=0$. In between these two special instances, a wide range of possible degrees of physical structure exist, degrees that should be reflected in the features of the underlying probability distribution. In the case of a PRNG, the 'ideal' values are $H[P]=1$ and $C[P]=0$.

As pointed out above, P itself is not a uniquely defined object, and several approaches have been employed in the literature so as to 'extract' P from the given time series. Just to mention some frequently used extraction procedures: (i) time series histogram (Martin 2004), (ii) binary symbolic dynamics (Mischaikow *et al.* 1999), (iii) Fourier analysis (Powell & Percival 1979), (iv) wavelet transform (Blanco *et al.* 1998; Rosso *et al.* 2001), (v) partition entropies (Ebeling & Steuer 2001), (vi) permutation entropy (Bandt & Pompe 2002; Keller & Sinn 2005), (vii) discrete entropies (Amigó *et al.* 2007), etc. There is ample liberty to choose among them. De Micco *et al.* (2008) proposed two probability distributions as relevant for testing the uniformity of $\mu(x)$ and the mixing constant: (i) a P based on time-series' histograms and (ii) a P based on ordinal patterns (permutation ordering) that derives from using the Bandt–Pompe method (Bandt & Pompe 2002).

For extracting P via the histogram, one divides the interval $[0, 1]$ into a finite number N_{bin} of non-overlapping subintervals A_i : $[0, 1] = \cup_{i=1}^{N_{\text{bin}}} A_i$ and $A_i \cap A_j = \emptyset \forall i \neq j$. Note that N in equation (3.1) is equal to N_{bin} . Of course, in this approach, the temporal order of the time series plays no role at all. In this paper, the quantifiers obtained via the ensuing PDF are called $H^{(\text{hist})}$ and $C^{(\text{hist})}$. Let us stress that for time series within a finite alphabet, it is relevant to consider an optimal value of N_{bin} (e.g. De Micco *et al.* 2008).

In extracting P by recourse to the Bandt–Pompe method, the resulting probability distribution P is based on the details of the attractor-reconstruction procedure. *Causal information* is, consequently, duly incorporated into the construction process that yields P . In this paper, the quantifiers obtained via the ensuing PDF are called $H^{(\text{BP})}$ and $C^{(\text{BP})}$. A notable Bandt–Pompe result consists in getting a clear improvement in the quality of information theory-based quantifiers (Larrondo *et al.* 2005, 2006; Kowalski *et al.* 2007; Rosso *et al.* 2007a,b, 2008; Zunino *et al.* 2007, 2008).

The extracting procedure is as follows. For the time series $\{x_t : t = 1, \dots, M\}$ and an embedding dimension $D > 1$, one looks for 'ordinal patterns' of order D (Bandt & Pompe 2002; Keller & Lauffer 2003; Keller & Sinn 2005) generated by

$$(s) \mapsto (x_{s-(D-1)}, x_{s-(D-2)}, \dots, x_{s-1}, x_s), \quad (3.6)$$

which assign to each 'time s ' a D -dimensional vector of values pertaining to the times $s, s-1, \dots, s-(D-1)$. Clearly, the greater the D -value, the more is the information on 'the past' incorporated into these vectors. By the 'ordinal pattern' related to the time (s) , we mean the permutation $\pi = (r_0, r_1, \dots, r_{D-1})$

of $(0, 1, \dots, D - 1)$ defined by

$$x_{s-r_{D-1}} \leq x_{s-r_{D-2}} \leq \dots \leq x_{s-r_1} \leq x_{s-r_0}. \quad (3.7)$$

In order to get a unique result, we consider that $r_i < r_{i-1}$ if $x_{s-r_i} = x_{s-r_{i-1}}$. Thus, for all the $D!$ possible permutations π of order D , the probability distribution $P = \{p(\pi)\}$ is defined by

$$p(\pi) = \frac{\sharp\{s | s \leq M - D + 1; (s) \text{ has type } \pi\}}{M - D + 1}. \quad (3.8)$$

In the last expression, the symbol \sharp stands for ‘number’.

The advantages of the Bandt–Pompe method reside in (i) its simplicity, (ii) the associated extremely fast calculation process, (iii) its robustness in the presence of observational and dynamical noise, and (iv) its invariance with respect to nonlinear monotonous transformations. The Bandt–Pompe methodology is not restricted to a time-series representative of low-dimensional dynamical systems but can be applied to any type of time series (regular, chaotic, noisy or reality based), with a very weak stationary assumption (for $k = D$, the probability for $x_t < x_{t+k}$ should not depend on t ; Bandt & Pompe 2002). One also assumes that enough data are available for a correct attractor reconstruction. Of course, the embedding dimension D plays an important role in the evaluation of the appropriate probability distribution because D determines the number of accessible states $D!$. Also, it conditions the minimum acceptable length $M \gg D!$ of the time series that one needs in order to work with a reliable statistics. In relation to this last point, Bandt and Pompe suggest, for practical purposes, working with $3 \leq D \leq 7$ with a time lag $\tau = 1$. This is what we do here (in the present work $D = 6$ is used).

(b) Quantifiers based on recurrence plots

Recurrence plots were introduced by Eckmann *et al.* (1987) so as to visualize the recurrence of states during phase space evolution. The recurrence plot is a two-dimensional representation in which both axes are time ones. The recurrence of a state appearing at two given times t_i, t_j is pictured in the two-dimensional graph by means of either black or white dots, where a black dot signals a recurrence. Of course only periodic continuous systems will have exact recurrences. In any other case, one detects only approximate recurrences, up to an error ϵ . The so-called recurrence function can be mathematically expressed as

$$\mathbf{R}(i, j) = \Theta(\epsilon - \|\vec{x}(i) - \vec{x}(j)\|), \quad (3.9)$$

with $\vec{x}(i) \in \mathfrak{R}^m$ and $i, j = 1, \dots, N$, N being the number of discrete states $\vec{x}(i)$ considered, $\|\cdot\|$ is a norm and $\Theta(\cdot)$ is the Heaviside step function.

In the particular case of the PRNGs analysed in this paper, only $1D$ series are considered, but the recurrence function idea can be extended to D -dimensional phase spaces or even to suitably reconstructed embedding phase spaces. Of course, the visual impact produced by the recurrence plot is insufficient to compare the quality of different PRNGs because of the ‘small-scale’ structures that may be present in our scenario. Several kinds of measures have been defined to quantify these small-scale structures (Marwan *et al.* 2007), each measure being a functional

of the recurrence function (equation (3.9)). In this paper, two kinds of recurrence plot measures are considered.

- (i) *Measures based on the recurrence density (measured by the number of points in the recurrence plot)*. In this paper, we use the *recurrence rate* (RR), given by

$$\text{RR}(\varepsilon) = \frac{1}{N^2} \sum_{i,j=1}^N \mathbf{R}_{ij}(\varepsilon). \quad (3.10)$$

Note that in the limit $N \rightarrow \infty$, RR is the probability that a state recurs to its ε -neighbourhood in phase space. For PRNGs, the ideal value would be $\text{RR} = 0$. But in practice, if no points are to be found in the recurrence plot, a larger ε must be adopted in order that the quantifier may make sense.

- (ii) *Diagonal measures*. These are measures related to the histogram $P(\varepsilon, l)$ of diagonal line lengths, given by

$$P(\varepsilon, l) = \sum_{i,j=1}^N [1 - \mathbf{R}_{i-1,j-1}(\varepsilon)] [1 - \mathbf{R}_{i+l,j+l}(\varepsilon)] \cdot \prod_{k=0}^{l-1} \mathbf{R}_{i+k,j+k}(\varepsilon). \quad (3.11)$$

Processes with uncorrelated or weakly correlated behaviour originate no (or just very short) diagonals, whereas deterministic processes give rise to ‘long’ diagonals and smaller amount of single, isolated recurrence points. In this paper, three measures based on the statistics of diagonal lines are considered.

- (a) *The deterministic quantifier DET*. The ratio of recurrence points that form diagonal structures of at least length l_{\min} to all recurrence points

$$\text{DET} = \frac{\sum_{l=l_{\min}}^N l \cdot P(\varepsilon, l)}{\sum_{l=1}^N l \cdot P(\varepsilon, l)}. \quad (3.12)$$

- (b) *The average diagonal line length L* given by

$$L = \frac{\sum_{l=l_{\min}}^N l \cdot P(\varepsilon, l)}{\sum_{l=l_{\min}}^N P(\varepsilon, l)}. \quad (3.13)$$

- (c) *The entropy ENTR* given by

$$\text{ENTR} = - \sum_{l=l_{\min}}^N P(\varepsilon, l) \ln P(\varepsilon, l). \quad (3.14)$$

(c) Quantifiers based on intrinsic computation

We consider in this paper two quantifiers introduced in Feldman *et al.* (2008), i.e. the *entropy rate* h_μ and the *entropy excess* \mathbf{E} . They are defined for time series with a finite alphabet \mathcal{A} , which is not a limitation because the x_i 's may be thought of as real numbers only in analytical studies. In any practical case, they are in fact *floating point numbers* and, consequently, they have only a finite number of allowed A -values. A subsequence $s^L = \{x_i, x_{i+1}, \dots, x_{i+L}\}$ is called an L -block.

Let $P(s^L)$ denote the probability of a particular L -block. Then the block entropy $H(L)$ is

$$H(L) \equiv - \sum_{s^L} P(s^L) \log_2 P(s^L). \quad (3.15)$$

The sum runs over all possible blocks of length $L > 0$ and $H(0) \equiv 0$ by definition. For stationary processes and sufficiently large L , $H(L) \sim L$. On the other hand, the entropy rate h_μ is defined as

$$h_\mu(L) = \frac{H(L)}{L} \quad h_\mu = \lim_{L \rightarrow \infty} h_\mu(L). \quad (3.16)$$

The entropy rate is also known as the metric entropy in dynamical systems theory and it is equivalent to the *thermodynamic entropy density* familiar from equilibrium statistical mechanics (Feldman *et al.* 2008). The entropy rate provides a reliable and well-understood measure of the randomness or disorder intrinsic to a process. However, this tells us little about the underlying system's organization, structure or correlations. A measure of the system's organization may be obtained by looking at the manner in which $h_\mu(L)$ converges to its asymptotic value h_μ . When only observations over length L -blocks are taken into account, a process appears to have an entropy rate of $h_\mu(L)$ larger than the asymptotic value of h_μ . As a result, the process seems to be of a more random nature than it really is by the 'excess' amount of $h_\mu(L) - h_\mu$ bits. Summing these entropy over estimates over L , one obtains the excess entropy (Crutchfield & Packard 1983)

$$\mathbf{E} \equiv \sum_{L=1}^{\infty} [h_\mu(L) - h_\mu]. \quad (3.17)$$

(d) *Expected behaviour for PRNG*

Summing up, the quantifiers to be compared here are: $H^{(\text{hist})}$, $C^{(\text{hist})}$, $H^{(\text{BP})}$, $C^{(\text{BP})}$, RR, DET, ENTR, L , h_μ and \mathbf{E} . These quantities should tell us how good our PRNG is as compared to the ideal condition $\mu(x) = \text{const.}$, $r_{\text{mix}} = 0$. $H^{(\text{hist})}$ is the natural quantifier to measure a non-constant $\mu(x)$, with value 1 for the ideal PRNG. It does not depend on the order of appearance of a given time-series event, but only on the number of times such event takes place. As for $H^{(\text{hist})}$, it is not able to uncover any change in r_{mix} -values. Thus, to get a good representation plane, we ought to demand a quantifier that changes with r_{mix} and not with $\mu(x)$. To look for such a kind of quantifier, we must study $H^{(\text{hist})}$, $C^{(\text{hist})}$, $H^{(\text{BP})}$, $C^{(\text{BP})}$, RR, DET, ENTR, L , h_μ and \mathbf{E} as functions of r_{mix} . A family of iterated maps f^d may be used to that end because they share the same invariant measure and r_{mix} is a decreasing function of d . The best quantifier for r_{mix} would be that which has maximal variation over the entire family of maps.

4. Application to logistic map and three-way Bernoulli map

In this section, we present results for the families of iterates of two chaotic well-known maps, the logistic map (LOG) and the three-way Bernoulli map (TWB), that have been selected, among other possibilities, because they are representative of two different classes of systems.

Table 1. Mixing constant (r_{mix} as a function of the iteration-order d for LOG and TWB).

d	TWB	LOG
1	0.56789	0.333333333
2	0.31848	0.111111111
3	0.13290	0.037037037
4	0.05788	0.012345679
5	0.03646	0.004115226
6	0.01791	0.001371742
7	0.01152	0.000457247
8	0.00515	0.000152416

(i) LOG is given by

$$x_{n+1} = 4x_n(1 - x_n), \quad (4.1)$$

and its natural invariant density can be exactly determined, being expressed in the fashion

$$\rho_{\text{inv}}(x) = \frac{1}{\pi\sqrt{x(1-x)}}. \quad (4.2)$$

LOG is paradigmatic because it is representative not only of maps with a quadratic maximum, but also emerges when the Lorenz procedure is applied to many continuous attractors with basins that may be approached with the Lorenz method via a one-dimensional map (like the Lorenz, Rossler and Colpits ones among others). A *non-uniform* natural invariant density is an important feature in this instance (Beck & Schlögl 1997). The ensuing r_{mix} -values are also displayed in table 1. They have been obtained by recourse to the transfer operator method, as described in Beck & Schlögl (1997).

(ii) TWB is given by

$$x_{n+1} = \begin{cases} 3x_n & \text{if } 0 \leq x_n \leq \frac{1}{3} \\ 3x_n - 1 & \text{if } \frac{1}{3} < x_n \leq \frac{2}{3} \\ 3x_n - 2 & \text{if } \frac{2}{3} < x_n \leq 1. \end{cases} \quad (4.3)$$

TWB is representative of the class of piecewise-linear maps as, for example, the four-way tailed shift map, the skew tent map, the three-way tailed shift map, etc. All these maps share a *uniform* natural invariant density (Beck & Schlögl 1997). The mixing constant r_{mix} of the whole family of maps f^d is given by $r_{\text{mix}}^d = (1/3)^d$ (table 1).

For the evaluation of the different quantifiers, we used files with $M = 50 \times 10^6$ floating point numbers. In the Band-Pompe approach, we consider $D = 6$ while

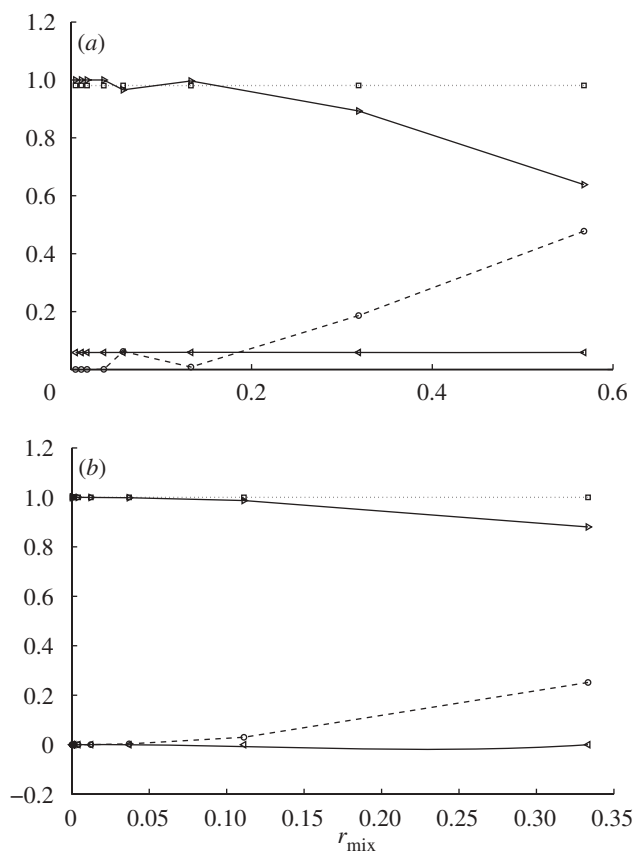


Figure 1. Information theory quantifiers as functions of r_{mix} : (a) LOG; (b) TWB. Squares, $H^{(\text{hist})}$; right triangles, $H^{(\text{BP})}$; left triangles, $C^{(\text{hist})}$; circles, $C^{(\text{BP})}$.

for histograms we have taken $N_{\text{bin}} = 2^{16}$. All recurrence plots measures depend on several parameters.

- (i) The dimension D_e of the embedding space. In this paper $D_e = 1$.
- (ii) ε , a parameter crucial so as to define just when recurrences occur. We adopted $\varepsilon = 1/(2^{16} - 1)$ corresponding to 16-bit numbers.
- (iii) l_{min} is the minimum length accepted for diagonal lines. $l_{\text{min}} = 2$ is used in this paper for all diagonal measures except for L ($l_{\text{min}} = 1$ is used for L).
- (iv) N is the number of values used for each realization. In this paper, values of RR, DET, ENTR and L are mean values over 10 surrogate series with $N = 10\,000$ data each.

Figures 1–3 illustrate the behaviour of all the quantifiers for the iterates of LOG (figures 1a, 2a and 3a) and the iterates of TWB (figures 1b, 2b and 3b). These figures show that the following quantifiers are the ones usable for measuring r_{mix} : $C^{(\text{BP})}$, DET, ENTR and L . On the other hand, the following quantifiers depend on the invariant density, but they do not depend on r_{mix} : $H^{(\text{hist})}$, $C^{(\text{hist})}$ and RR.

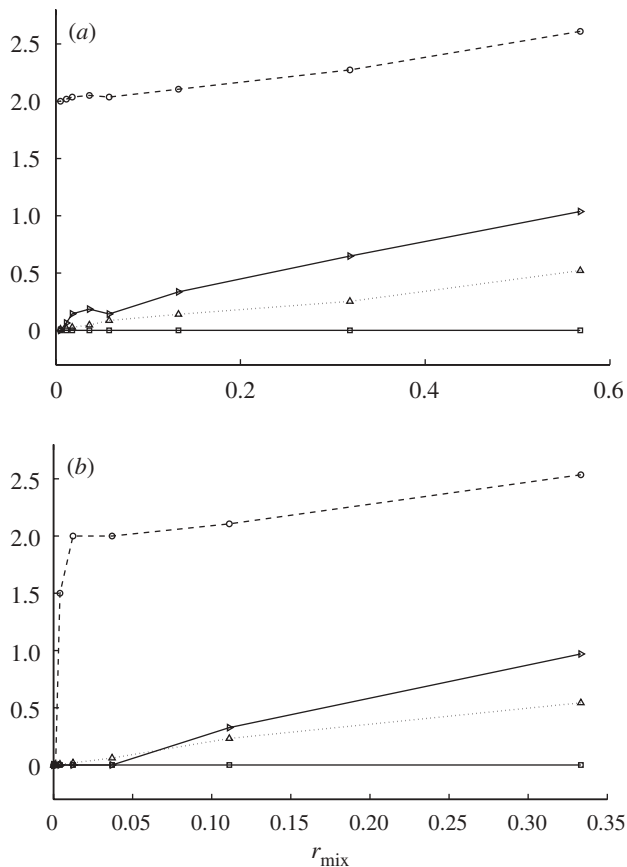


Figure 2. Recurrence plots quantifiers as functions of r_{mix} : (a) LOG; (b) TWB. Squares, RR; circles, L; up triangles, DET; right triangles, ENTR.

Intrinsic computation quantifiers display a completely different behaviour for LOG and for TWB. In LOG, both quantifiers have no dependence with r_{mix} , but in TWB, h_{μ} decreases as r_{mix} increases while \mathbf{E} is an increasing function of r_{mix} . Thus, these parameters do not seem to be convenient ones.

Comparing LOG with TWB by recourse to these parameters shows that TWB is slightly better than LOG. The problem with TWB and with other piecewise-linear maps is they are not realistic enough and that their implementation is more involved than for other maps like LOG.

As an application of the above quantifiers, we study two usual randomization procedures by means of the representation plane depicted in figure 4, employing a quantifier depending on $\mu(x)$ as x -axis ($H^{(\text{hist})}$ is selected) and a quantifier depending on r_{mix} as y -axis (DET is selected). The first procedure is skipping and the second one is discretization (De Micco *et al.* 2008). Skipping has been used as a randomization procedure for piecewise-linear maps. The representation plane evidences the fact that this procedure is better than discretization because these maps already have the ideal invariant measure (they have $H^{(\text{hist})} = 1$) and only r_{mix}

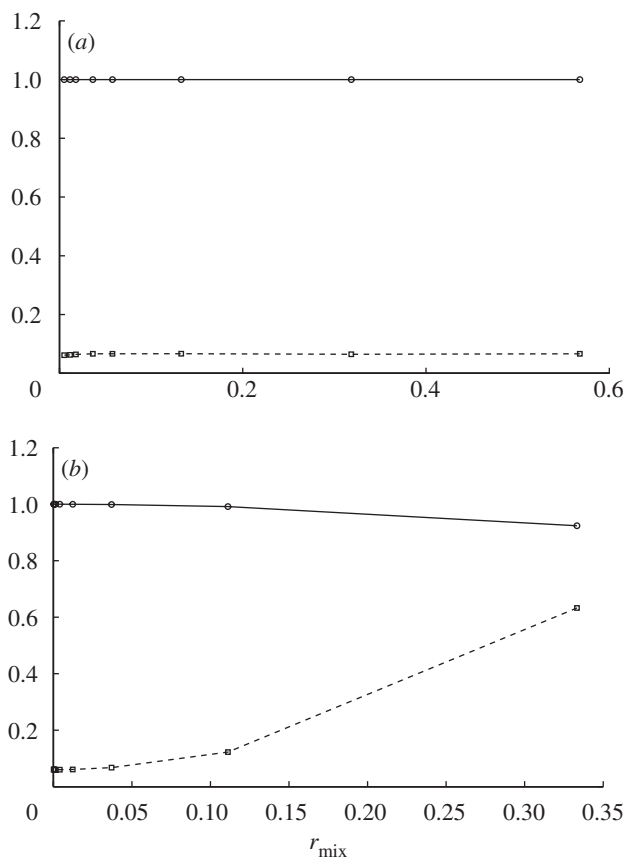


Figure 3. Intrinsic computation quantifiers as functions of r_{mix} : (a) LOG; (b) TWB. Squares, \mathbf{E} ; circles, h_{μ} .

must be diminished to get the ideal PRNG. Figure 4b shows that discretization is a better procedure for LOG because the ideal point $[1, 0]$ is not reached by skipping.

5. Conclusions

In summary, we were able to show here the following.

- (i) Two classes of quantifiers are required for the evaluation of the quality of a PRNG: (i) *quantifiers depending on r_{mix} only (and not on $\mu(x)$)*, like $H^{(\text{hist})}$, $C^{(\text{hist})}$ and RR and (ii) *quantifiers depending on $\mu(x)$ only (and not on r_{mix})*, as $C^{(\text{BP})}$, $H^{(\text{BP})}$, DET , L and $ENTR$.
- (ii) Intrinsic computation quantifiers are dependent on both $\mu(x)$ and r_{mix} and then they are not convenient for PRNG analysis with our methodology.
- (iii) Representation planes with one quantifier of each class as coordinate axis allow for different chaotic PRNGs to be compared with each other so as to determine the best one. Furthermore, these representation planes permit one to judiciously select the best randomizing procedure.

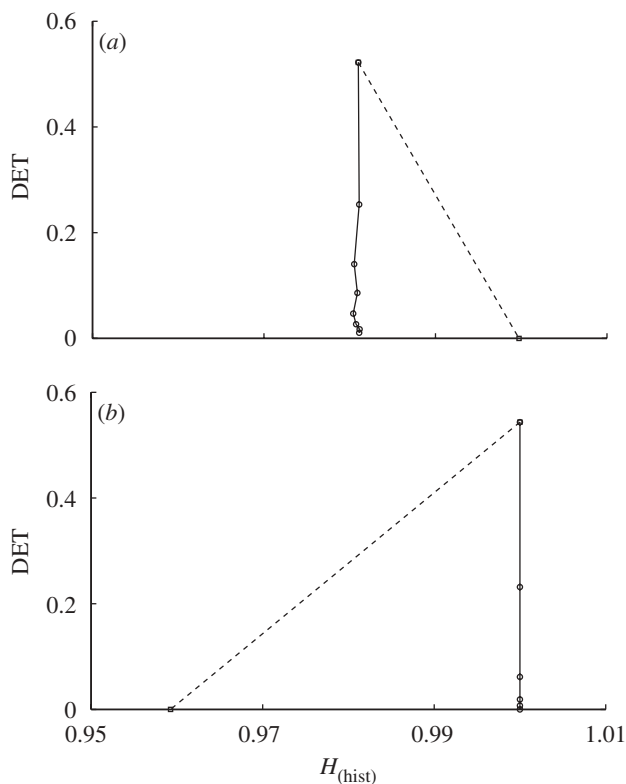


Figure 4. DET as a function of $H^{(\text{hist})}$, as evaluated in De Micco *et al.* (2008), for both randomization procedures applied to: (a) LOG; (b) TWB. Squares and dashed lines, discretization; circles and solid lines, skipping.

Our present results are consistent with those of previous works (Gonzalez *et al.* 2005; Larrondo *et al.* 2005; De Micco *et al.* 2008).

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