# A System for Offline Cursive Handwritten Word 

# Recognition 


#### Abstract

by Marlon Abayan Submitted to the Department of Electrical Engineering and Computer Science in Partial Fulfillment of the Requirements for the Degrees of Bachelor of Science in Computer Science and Engineering and Master of Engineering in Electrical Engineering and Computer Science at the

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#### Abstract

Cursive handwriting recognition is a difficult problem because of large variations in handwritten words as well as overlaps and interconnections between neighboring characters. In this thesis we introduce a new approach to this problem, called Hidden Markov Model with Multiple Observation Sequences (HMMMOS). A preprocessor extracts each word from a scanned-in document image and divides it into segments. A Neural Network (NN) classifier then finds the likelihoods of each possible character class given the segments and combinations of segments. These likelihoods, along with statistics computed from a lexicon, are used as input to a dynamic programming algorithm which recognizes the entire word. The dynamic programming algorithm can be viewed as a modified Viterbi algorithm for a Hidden Markov Model (HMM). Three types of Neural Networks are tried: a recurrent network, a Multilayer Perceptron (MLP), and a Hierarchical Mixture of Experts (HME) [Jordan \& Jacobs 1994]. As an extension, the Viterbi algorithm for the HMM is implemented for a multiprocessor environment to speed up recognition.


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## Chapter 1

## Introduction

### 1.1 Problem Definition

This thesis proposes a solution to the off-line cursive handwritten word recognition problem. Handwritten word recognition (HWR) is the problem of interpreting a word image by assigning to it a particular sequence of characters called a "word". The subproblems range from recognizing block-printed and disconnected characters to identifying cursively handwritten words. HWR can be done on-line or off-line, as differentiated in [Seiler, Schenkel, \& Eggimann 1996]. When the recognition is done on-line, a user typically writes a word onto a computer interface. Temporal information, as well as the word image, is available to the recognition system. When it is done off-line, however, only the handwritten word image is available. Hence, off-line handwriting recognition is more difficult.

Examples of real world applications of off-line cursive HWR include personal check processing, mail sorting, and automatic data entry of business forms.


Figure 1.1: Overall System Specification

### 1.2 Background Literature

The strategies for off-line cursive handwritten word recognition can be roughly classified
into three categories. In the first category, the word is segmented into several characters and character recognition techniques are applied to each segment [Casey and Nagy, 1982]. This method heavily depends on the accuracy of the segmentation points found. Unfortunately, such an accurate segmentation technique is not yet available. In the second category, whole words are recognized without doing any kind of formal segmentation [Farag 1979] [Hull \& Srihari 1986]. The global shape is analyzed to hypothesize the words. This is effective for small lexicons such as in check processing. The difference between the global shapes of "one", "two", and "dollar" can easily be captured by a simple classifier. As the lexicon grows, however, the differences between words become more minute and the classifier needs to be much more complex. Usually this strategy does not work for medium-sized and large lexicons. The third category is a compromise solution between these two schemes. It does a loose segmentation to find a number of potential segmentation points [Bozinovic \& Srihari 1986] [He, Chen, Kundu 1996]. The final segmentation and word identity are determined later in the recognition stage with the help of a lexicon. Figure 1.2 illustrates this two step process.


Figure 1.2: Two-step approach

The new Hidden Markov Model (HMM) approach to HWR described in this paper belongs to the third category.

### 1.3 Hidden Markov Model (HMM) Framework

Like most real world processes, handwriting produces observable outputs which can be characterized as sequential signals. These signals can be successfully coded by a stochastic model such as a Hidden Markov Model. It is quite natural to model handwriting as an embedded Markov chain in an HMM. For example, a word can be regarded as a Markov chain of individual characters.

An HMM has an underlying Markov process that is "hidden" or not directly observable, i.e. it does not produce the observation sequences directly. Another set of stochastic processes produces the sequence of observed symbols and stochastically relate the symbols to the "hidden" Markov states. They are called observation probability distributions.

HMMs are often called doubly stochastic processes because there are two sets of processes. They are more powerful than Markov models which allow one less degree of freedom. In handwriting, the writer translates the language to the written text with added ambiguities, e.g., variation in writing style, ambiguity caused by connected and overlapping letters, etc. These ambiguities can be accommodated by the observation probability distributions.

A particular difficulty in the use of HMMs for HWR is the following problem. The potential segmentation points found in the loose segmentation phase are a superset of the points that would partition the image into the underlying characters. It is therefore necessary to identify which subsegments, when joined, form a single character. [Chen et. al. 1995] successfully applied a Variable Duration Hidden Markov Model (VDHMM) to solve this problem. Each state is actually a "small HMM" which itself has four states to account for up to four segments. The small HMMs embody the variability in the duration (the number of segments) of each state. In order to run a Viterbi algorithm for them, there is a need to estimate duration statistics.

### 1.4 Variable Sequence Length HMM

[He et. al. 1996] discusses a similar approach which does not involve computing duration statistics, called Variable Sequence Length HMM (VSLHMM). Instead, it explores the possible combinations of character subsegments using a modified dynamic programming Viterbi algorithm. However, it is not a pure dynamic programming algorithm because it has a greedy component which arises from the following assumption:

> "If one of the 4 block images consisting of either 1 or 2 or 3 or 4 consecutive segments (starting from the same segment) is an actual character image of a [character state], then if we compare the four images with that letter in the feature space, the actual character image always has the best match."

The characters are oversegmented into at most four segments. However, for consistency with our work, assume that the oversegmented characters are composed of only up to three segments.

The next section explains how the assumption can break and the consequences of that happening. This section steps through how the algorithm works on a sample image (Figure 1.3).


Figure 1.3: Sample oversegmented word

Figure 1.4 shows a dynamic programming table used to find an optimal word, i.e. state sequence, for Figure 1.3. Each row corresponds to a character state; each column, $\mathrm{k}=1,2$, $\ldots, 8$, corresponds to the $\mathrm{k}^{\text {th }}$ successive letter in the state sequence. The entries of the table
consist of a pair: the starting segment number, and the number of segments used up by the character. Two paths through the table are shown, one dashed and the other solid. The solid is the correct state sequence: ' $e$ ', ' $x$ ', ' $a$ ', ' $m$ ', ' $p$ ', ' 1 ', ' $e$ '. It starts with an ' $e$ ' because the first column entry along the path is on the ' $e$ ' row. The entry, $\langle 1,1\rangle$, means that the starting segment is 1 and ' $e$ ' takes up exactly 1 segment. The second state is ' $x$ ' because the next entry along the path is on row ' $x$ '. Similarly, the third state is ' $a$ '. Column 4 shows that the fourth state is ' $m$ '. This character starts at segment 4 and takes up 3 segments (see Figure 1.3). The remaining states are ' p ', ' l ', and ' e ', each taking up one segment.

The dashed path is ' $v$ ', ' $a$ ', ' $v$ ', ' $v$ ', ' $p$ ', ' $b$ '. It has six states and therefore uses up six columns. The first state, ' $v$ ', takes up the first two segments. The second state, ' $a$ ', uses up only the third segment, and so on along the path. Notice that both paths cross at row ' p ' and column 5.


Figure 1.4: Modified Viterbi

The algorithm explores many possible paths through the table and the one with the highest probability is chosen. The state sequence $\mathrm{Q}=\left\langle\pi, \mathrm{q}_{1}, \mathrm{q}_{2}, \ldots, \mathrm{q}_{\mathrm{T}}>\right.$, implied by a path through the table, is mapped to a probability by the expression below.

$$
a_{\pi, q_{1}} P\left(o_{1} \mid q_{1}\right) a_{q_{1}, q_{2}} \ldots a_{q_{T-1}, q_{T}} P\left(o_{T} \mid q_{T}\right)
$$

Each $\mathrm{o}_{\mathrm{i}}$ is the combination of adjacent segments derived from the table entry along the path at column i. Let us define any combination of up to three adjacent segments to be a chunk. Each $\mathrm{o}_{\mathrm{i}}$ is therefore a chunk.

The algorithm starts by filling in the first column. All the rows (states) have to begin at the first segment so the first element of each table entry on the first column is always 1. The second element is determined by finding the number of segments which will optimize the probability $\mathrm{P}(\mathrm{olq})$ where q is the state and o is the chunk composed of the segments. This step arises directly from the assumption mentioned above. The number of segments "consumed" by a state at a given column is greedily set to the one that maximizes the probability at that point in the algorithm.

The second column is filled by finding the state in the first column, $\mathrm{i}^{*}$, that maximizes the path probability up to that point,

$$
i^{*}=\underset{i}{\arg \max } \quad a_{\pi, q_{i}} P\left(o_{1} \mid q_{i}\right) a_{q_{i} q_{j}} P\left(o_{2} \mid q_{j}\right)
$$

where i is a previous state and j is the current state. Back-pointers to $\mathrm{i}^{*}$ for each j are also stored. The chunks $\mathrm{o}_{1}$ and $\mathrm{o}_{2}$ are selected by maximizing $\mathrm{P}\left(\mathrm{o}_{1} \mid \mathrm{q}_{\mathrm{j}}\right)$ and $\mathrm{P}\left(\mathrm{o}_{2} \mid \mathrm{q}_{j}\right)$, respectively, for various combinations of segments. The algorithm proceeds in the same fashion for subsequent columns until all paths have used up all the segments. Then, the path with the highest probability is selected.

### 1.5 HMM with Multiple Observation Sequences

This thesis introduces an HMM architecture coined HMM with Multiple Observation Sequences (HMMMOS). It is based on VSLHMM, but relaxes the assumption stated in section 1.4 because it frequently breaks in practice and problems arise.

For example, if segment 2 in Figure 1.3 is always merged with segment 1 for all the states with high probability in column 1 , it will never be recognized as a separate character ' $x$ '. That is, if all table entries in the first column are of the form $\langle 1,2\rangle$ or $\langle 1,3\rangle$ then no path in the table will ever have a separate segment 2 . For state ' $e$ ' the assumption that the best chunk will have the highest probability is violated. Instead, the chunk composed of segments 1 and 2 , or 1 and 2 and 3 , have the highest probability.

Similarly, if segment 3 is merged with segment 4 for state 'a' we can never separate them. The resulting misalignment eats into the ' $m$ ' character and causes more errors down the line.

These problems are solved by HMMMOS by having a more general dynamic programming Viterbi algorithm which gets rid of the greedy component. The algorithm is discussed in detail in Chapter 3.

Although HMMs and HMMMOS are very similar in that they have Viterbi-like procedures for assigning probabilities to candidates, there are some major differences. An HMMMOS is not trained using the canonical HMM training algorithm [Rabiner 1989]. Instead, the transition probabilities are estimated directly from a lexicon and the observation probability distributions are provided by a neural classifier which is trained separately. Furthermore, the units of inputs, the word segments, do not generally correspond to a single state: groups of inputs do. This causes the Viterbi algorithm to be more complicated.

In order to characterize HMMMOS more explicitly than as an extended HMM, the next section compares it with Stochastic Context Free Grammars (SCFG), and the associated training and recognition algorithm, the Inside-Outside Algorithm.

### 1.6 Stochastic Context Free Grammars and Inside-Outside Algorithm

HMMMOS and SCFGs are both stochastic models for sequential data. Their main similarity is in the training and the fact that segments are recombined into chunks, which looks like combining nonterminals in a grammar production.

A basic formulation of stochastic context-free grammars is given by [Fu 1976]. A con-text-free language is one in which a grammar in the Chomsky normal form [Chomsky 1959] can be found. I.e., productions can only be of the form

$$
\begin{gathered}
i \rightarrow j k \\
\quad \text { or } \\
i \rightarrow m
\end{gathered}
$$

where the symbols $\mathrm{i}, \mathrm{j}$, and k are nonterminals and m is a terminal symbol. A stochastic context-free grammar is a grammar for a context-free language in which all productions are associated with a probability. The probabilities for the grammar rules for a context-free grammar G are stored in matrices A and B with elements

$$
\begin{gathered}
a[i, j, k]=P(i \rightarrow j k \mid G) \\
b[i, m]=P(i \rightarrow m \mid G)
\end{gathered}
$$

representing the probability of the productions of the forms $\mathrm{i}->\mathrm{jk}$ and $\mathrm{i}->\mathrm{m}$, respectively. Therefore, the parameters stored in the A matrix represent the hidden process while the parameters stored in the B matrix represent the observable process.

The Inside-Outside algorithm [Lari \& Young 1991] is used to train parameters from sample data and recognize test sentences given the grammar. The algorithm defines inner (e) and outer (f) probability distributions as follows,

$$
\begin{gathered}
e(s, t, i)=P(i \Rightarrow O(s), \ldots, O(t) \mid G) \\
f(s, t, i)=P(S \Rightarrow O(1), \ldots, O(s-1), i, O(t+1), \ldots, O(T) \mid G)
\end{gathered}
$$

The symbol "=>" means that the symbol on the left hand side eventually expands into the right hand side using zero or more productions. $S$ is the start state of the grammar. The algorithm proceeds by computing the inner probabilities in a bottom-up fashion and the outer-probabilities in a top-down fashion.

An HMMMOS can be interpreted as an SCFG in the following way. All the states of the HMMMOS, the characters, are nonterminals in the context-free grammar. The image segments are terminal symbols. State sequences (words and subwords) are also nonterminals. Words are formed by expanding word nonterminals into pairs of nonterminals in productions of the form $\mathrm{i}->\mathrm{jk}$, until the expression is expressed purely in terms of state nonterminals. The state nonterminals are then expanded into productions involving the terminal symbols, the segments. Here are some examples of productions for the state ' $w$ ':

$$
\begin{gather*}
w \rightarrow w_{L} w_{R}  \tag{1}\\
w_{R} \rightarrow \alpha_{1}  \tag{2}\\
w_{L} \rightarrow w_{L_{L}} w_{L_{R}}  \tag{3}\\
w_{L} \rightarrow \alpha_{2}  \tag{4}\\
w \rightarrow \alpha_{3}  \tag{5}\\
w_{L_{L}} \rightarrow \alpha_{4}  \tag{6}\\
w_{L_{R}} \rightarrow \alpha_{5} \tag{7}
\end{gather*}
$$

The $\alpha_{i}$ are nonterminal symbols corresponding to specific word segments.
There are many ways a letter, in this case ' $w$ ', can be formed. Each way has a different probability, with some of them being equal to zero. In production (1) the character is bro-
ken up into two nonterminals. The left nonterminal can be a nonterminal segment (4) or a combination of adjacent nonterminal segments (3), (6), and (7). The right nonterminal (1) can only be a segment (2). This makes the grammar for ' $w$ ' left recursive, though it need not be. Finally, the character could be formed by a single segment (5). Each of (2), (4), (5), (6), (7) actually represent a very large number of productions because the $\alpha_{i}$ are arbitrary bitmaps. If the size and resolution of the bitmaps are unconstrained, they correspond to an infinite number of productions. The probabilities for this class of productions are in the matrix B discussed above.

Using our segmentation algorithm, the character ' w' would most often be characterized by (1) and (3) because ' $w$ ' has three ascending spikes. The character ' $n$ ', on the other hand, would most often be characterized by (1) and (4) because it has two descending spikes, and ' $a$ ' by (5) because it has a roundish self-contained shape.

The main difference between the HMMMOS and SCFG arises when probabilities are assigned to the productions. HMMMOS does not assign probabilities to (2), (4), (6), and (7). Instead, probabilities are assigned directly to (1) and (5) based on the combined bitmap formed from all the segments involved. Hence, in contrast to SCFG, the intra-state productions of the form $\mathrm{i}->\mathrm{jk}$ are not independent of the productions of the form $\mathrm{i}->\mathrm{m}$. The Viterbi algorithm for HMMMOS uses a bottom-up dynamic programming algorithm to consider different combinations of adjacent segments into chunks which form the right hand side of (1) and (5).

The transition probabilities in HMMMOS is the same as the production probabilities for contiguous characters in the context-free grammar. The observation probabilities are exactly the elements of matrix B. Matrix A in the SCFG is just the transition probabilities, along with the probabilities of productions of single characters comprised of chunks of
segments just described. Matrix B is the set of observation probabilities and a set of undefined and unused probabilities for productions of the form (2), (4), (6), and (7).

Training-wise, HMMMOS and SCFG are similar in that the A and B matrices are determined separately but use the same data. The exception, of course, is the probabilities in productions (1) and (3) in matrix A can only be determined after matrix $B$ has been determined.

The HMMMOS Viterbi algorithm is similar to the Inside-Outside algorithm [Lari \& Young 1991] in that it computes the probability in a dynamic programming manner. The Inside-Outside algorithm does a bottom-up for inner probabilities and top-down for outer probabilities. The Viterbi algorithm only does a bottom-up computation as described in detail in Chapter 3.

### 1.7 Organization of this thesis

An overview of the software system, the main modules, and the contributions of this thesis is given in Chapter 2. The underlying idea for this thesis, the HMM with Variable Observation Sequences, is described in Chapter 3. Chapter 4 explains how the different probabilities in the model are estimated. Chapter 5 shows a verification algorithm that ranks candidate words for a given oversegmented word image, and Chapter 6 discusses a neural classifier used to reject bad inputs and improve the recognition rate. Then, Chapter 7 extends the work by introducing a parallel implementation of a simplified version of the system. Each chapter gives the results of the algorithms it presents. Finally, Chapter 8 concludes and suggests future work.

## Chapter 2

## System Overview

### 2.1 Architecture

This chapter outlines the software architecture and gives an overview of the main modules. As is typical among supervised learning systems, there is a training phase and a testing (or recognition) phase. The data, which come in the form of word images, is divided into two sets: the training set, for training the neural network, and the test set, for evaluating the performance of the system. Figure 2.1 shows a data flow diagram.

The majority of the code is written in C for an RS/6000 machine running AIX.


Figure 2.1: Software Infrastructure

### 2.2 Training

The purpose of the training phase is to train the neural network that generates the observation (character bitmap) probabilities and estimate state transition probabilities.

### 2.2.1 Oversegmentation

After word images are extracted from the input images ${ }^{1}$ they are oversegmented. Continuous handwriting is cut into pieces by specifying the beginning and ending of the stroke. Pairs of points mark the top and the bottom pixels of these locations. Therefore, each segment has up to four cutting points. Segments that are at the beginning and end of a connected component only have two cutting points.

Each word is segmented into letters and parts of letters. A particular letter may have one, two, or up to three segments. All the letters in the word can are made up of segments or unions of adjacent segments. Figure 2.2 shows some examples of actual oversegmented word images.



Wocestr


Figure 2.2: Oversegmentation examples.

1. In general, input images contain multiple words and even multiple lines of text.

### 2.2.2 Truthing

From segmented words, chunks that correspond to whole characters are manually chosen. They are manually labeled with the correct character states in a process called truthing. Only the truthed chunks are needed for training.

### 2.2.3 Feature Extraction

A feature vector composed of 108 features is extracted from each truthed chunk. We use 88 contour direction features [Mohiuddin \& Mao 1994] and 20 cutting point features.

Contour direction features can be efficiently extracted by scanning the normalized $24 \times 16$ image by a $2 \times 2$ mask to detect primitive features when the number of black pixels in the mask is between one and three (neither all white nor all black). Histograms are computed by summing the number of primitive features detected along each of the horizontal, vertical, and two diagonal directions. Therefore, each detected direction is assigned four times according to its position in four projections. The vertical projection has four slices and the others have six slices ( 22 slices in total). This results in an 88-dimensional (22 slices x 4 directions) feature vector.

The relative positions of the cutting points in the bounding box of merged segments (previous segment with the current segment, the current segment with the next segment) are coded in 20 additional features, resulting in a 108-dimensional feature vector.
2.2.4 Neural Network character recognizer

The feature vectors and truth labels are used to train the neural classifier. The neural classifier provides the set of probability distributions that map observations ${ }^{1}$ to the "hidden" states ${ }^{2}$ of the HMM.

There are many character recognition techniques being used for off-line cursive HWR. The most popular ones are artificial neural networks. Among the different types of neural networks, we try three: a Multilayer Perceptron (MLP) [Rumelhart, Widrow, Lehr 1994], a Recurrent Network which is based on an MLP architecture, and a Hierarchical Mixture of Experts (HME) [Jordan \&Jacobs 1994].

HME is a space partitioning scheme which was motivated by the fact that it can perform function decomposition and so learn separate sub-tasks within a large task more quickly and perhaps more effectively than a fully interconnected network [Jacobs, Jordan \& Barto 1991]. The scheme has also been applied to the classification problem of speech recognition [Waterhouse 1993]. We hope to use it in our application in a similar manner to cope with high input variance.

### 2.2.5 Transition Probabilities

While the neural classifier is being trained, transition probabilities are estimated using the frequencies of transitions in the truth labels. The training data set is used as an implicit lexicon.

### 2.3 Recognition

In the recognition phase, test input or new text images are assigned ASCII representations.

1. These are segments or combinations of up to three adjacent segments, otherwise known as chunks.
2. The hidden states are actually character classes, e.g. ' $a$ ', ' $b$ ', in the model.

New images are segmented and extracted of features as described in the previous section. Then all the chunks are fed into the neural classifier. A dynamic programming Viterbi algorithm, described in Chapter 3, is then applied to the recognition results to find word hypotheses. These hypotheses are ranked according to probability using global context and lexicon constraints. Bad samples are discarded to improve recognition rates.

### 2.4 Software modules developed in this thesis

Some of the software modules needed in the system had already existed at the IBM Almaden Research Center. The new modules and that were constructed to complete the system, as well as extensions, are:

1. Two-level dynamic programming Viterbi algorithm for the HMM.
2. Neural Networks (attempts at HME and recurrent network, as well as training the MLP using only whole characters).
3. Transition and initial probability estimation.
4. Recognition of segments and combinations of segments in the recognition phase.
5. A verification algorithm that ranks candidate words.
6. Rejection of bad samples to increase recognition rates (uses an MLP).
7. Parallel implementation of simpler problem to explore ways to improve performance.

## Chapter 3

## HMM with Multiple Observation Sequences

### 3.1 HMM Framework

We propose to use the following Hidden Markov Model (HMM) framework. The model is a directed graph where the nodes are called states. A state, i , in the model corresponds to a character class. For example class 1 corresponds to the letter ' $a$ ', 2 to letter ' $b$ ', and so forth. The set of character classes $\{1,2, \ldots, \mathrm{~N}\}$ contains all English letters, numbers, and a subset of the punctuation marks. Every path through the model must start from a special state called $\pi$. A typical path looks like $\mathrm{Q}=<\pi, \mathrm{q}_{1}, \mathrm{q}_{2}, \ldots, \mathrm{q}_{\mathrm{T}}>$, where $1 \leq \mathrm{q}_{\mathrm{i}} \leq \mathrm{N}$. Paths can be interpreted as words where T is the number of letters in the word.


Figure 3.1: Hidden Markov Model

There are links from one state to another. Each link is associated with the probability of making a transition from the source ${ }^{1}$ state to the destination state. These probabilities are called transition probabilities. They allow us to assign probabilities to entire words or paths through the graph. Consider a word $\mathrm{W}=$ " $\mathrm{c}_{1} \ldots \mathrm{c}_{\mathrm{T}}$ " where $\mathrm{c}_{\mathrm{i}}$ are characters for $\mathrm{i}=1$,

[^0]..., T. Let $\mathrm{q}_{\mathrm{c}}$ represent the state in the HMM corresponding to the character c . The word W can be represented as a state sequence $Q$
$$
Q=\left\langle\pi, q_{c_{1}}, q_{c_{2}}, \ldots, q_{c_{T}}\right\rangle
$$
with probability
$$
P(Q)=a_{\pi, q_{1}} a_{q_{1}, q_{2}} \ldots a_{q_{T-1}, q_{T}}
$$

In keeping with the laws of probability, the sum of the transition probabilities of all the outgoing links from any node must be equal to one.

$$
\sum_{k=1}^{N} a_{q_{i} q_{k}}=1 \quad \forall(1 \leq i \leq N)
$$

A link with probability equal to zero is equivalent to having no link at all. For example, all probabilities $\mathrm{a}_{\mathrm{i}, \pi}$ are equal to zero since no state can transition back to the $\pi$ state.

The initial state probabilities, $\mathrm{a}_{\pi, \mathrm{i}}$, and transition probabilities, $\mathrm{a}_{\mathrm{ij}}$, can be estimated from the training data or from a lexicon.

So far we have described a simple Markov Model. We extend this by introducing observation probability distributions $\mathrm{P}(o l i)$ at each state i . Here o is a chunk ${ }^{1}$ of a word image. The joint probability of a path and a perfectly segmented bitmap $O=<o_{1}, o_{2}, \ldots$, $\mathrm{o}_{\mathrm{T}}>$ is

$$
P(O, Q)=P(Q) P(O \mid Q)=\left(a_{\pi, q_{1}} a_{q_{1}, q_{2}} \ldots a_{q_{T-1}, q_{T}}\right)\left[P\left(o_{1} \mid q_{1}\right) P\left(o_{2} \mid q_{2}\right) \ldots P\left(o_{T} \mid q_{T}\right)\right]
$$

1. Recall that a chunk is a combination of up to three contiguous segments.

The distributions P (oli) come from a neural classifier that process feature vectors that have been extracted from the bitmap image of the corresponding chunk. An example of a neural classifier is a the Multilayer Perceptron (MLP).


Figure 3.2: Input/Output Relation for MLP used as Neural Classifier

The problem of finding the word represented by an image $S$ is therefore the exact same problem as identifying the state sequence, $\mathrm{Q}=<\pi, \mathrm{q}_{1}, \mathrm{q}_{2}, \ldots, \mathrm{q}_{\mathrm{T}}>$, that maximizes the joint probability $\mathrm{P}(\mathrm{Q}, \mathrm{S})$.

A word image $S$ can be represented by a sequence of segments, $S=\left\langle s_{1}, s_{2}, \ldots, s_{R}\right\rangle$. If we are lucky, S will be perfectly segmented. There is a one-to-one correspondence between the segments $\mathrm{s}_{\mathrm{i}}$ and the states $\mathrm{q}_{\mathrm{i}}$. So $\mathrm{R}=\mathrm{T}$. A Viterbi algorithm to find the underlying word, $\mathrm{Q}_{\mathrm{opt}}=<\pi, \mathrm{q}_{1}, \mathrm{q}_{2}, \ldots, \mathrm{q}_{\mathrm{T}}>$, that maximizes the likelihood of the word constructed by putting each segment in its own chunk, i.e.

$$
\begin{gathered}
Q_{\text {opt }}=\begin{array}{c}
\arg \max \\
q_{1}, q_{2}, \ldots, q_{T}
\end{array} P(Q, S) \\
={\underset{q}{1} 1}^{\arg \max } q_{2}, \ldots, q_{T}\left(a_{\pi, q_{1}} P\left(s_{1} \mid q_{1}\right) a_{q_{1} q_{2}} P\left(s_{2} \mid q_{2}\right) a_{q_{2} q_{3}} \ldots a_{q_{T-1} q_{T}} P\left(s_{T} \mid q_{T}\right)\right)
\end{gathered}
$$

This optimizes $\mathrm{P}(\mathrm{S}, \mathrm{Q})=\mathrm{P}(\mathrm{S} \mid \mathrm{Q}) \mathrm{P}(\mathrm{Q})$. It is equivalent to optimizing the a posteriori probability $\mathrm{P}(\mathrm{QIS})=\mathrm{P}(\mathrm{S}, \mathrm{Q}) / \mathrm{P}(\mathrm{S})$ since the word image is fixed.

In reality there is at present no segmentation algorithm that perfectly segments word images into the constituent characters. The word images in our system are usually oversegmented. At least one segment and at most three segments correspond to each character in virtually all cases.

To handle this general case, we incorporate in our Viterbi algorithm a search for the optimal grouping of the segments into characters. This search also uses a dynamic programming algorithm like the one used for the Viterbi algorithm. Therefore, the resulting modified Viterbi algorithm is a two-level dynamic programming algorithm.

### 3.2 Viterbi algorithm with multiple observation sequences.

This section describes the original Viterbi algorithm [Rabiner 1989] and the two-level dynamic programming "modified Viterbi algorithm". The latter is original work for this thesis.

### 3.2.1 Original Viterbi Algorithm

Let us consider the simple case where we would like to identify the single best word that would account for a perfectly segmented word image $S$. This word is the path Q through the HMM that maximizes $\mathrm{P}(\mathrm{SIQ})$. As shown above, this is the same path that maximizes $\mathrm{P}(\mathrm{Q} \mid S)$ after making some assumptions. To find this path the Viterbi algorithm [Viterbi 1967] [Rabiner 1989], which is a dynamic programming algorithm, is used.

To reiterate, we would like to find the best state sequence, $\mathrm{Q}=<\pi, \mathrm{q}_{1}, \mathrm{q}_{2}, \ldots, \mathrm{q}_{\mathrm{T}}>$, for the given observation sequence $S=\left\langle s_{1}, s_{2}, \ldots, s_{T}\right\rangle$.

Let us define the quantity

$$
\delta_{t}(i)=\begin{gathered}
\max \\
q_{1}, q_{2}, \ldots, q_{t-1}
\end{gathered} \quad P\left[q_{1}, q_{2}, \ldots, q_{t}=i, s_{1}, s_{2}, \ldots, s_{t}\right]
$$

I.e. $\delta_{t}(i)$ is the highest probability along a single path, at time $t$, which accounts for the first $t$ observations and ends in state $c_{i}$. By induction we have

$$
\delta_{t+1}(j)=\max _{i} \delta_{t}(i) a_{i j} P\left(s_{t+1} \mid j\right)
$$

To actually retrieve the state sequence, we need to keep track of the argument which maximized the equation above, for each $t$ and $j$. We do this via the array $\psi_{t}(j)$. The complete procedure for finding the best state sequence can be stated as follows:

1) Initialization:

$$
\begin{array}{cc}
\delta_{1}(i)=a_{\pi, i} P\left(s_{1} \mid c_{i}\right) \quad 1 \leq i \leq N \\
\psi_{1}(i)=0 &
\end{array}
$$

2) Recursion:

$$
\begin{array}{cc}
\delta_{t}(j) & =\begin{array}{c}
\max \\
1 \leq i \leq N
\end{array} \\
\psi_{t}(j) & =\begin{array}{l}
t-1 \\
\arg \max \\
1 \leq i \leq N
\end{array} \quad\left[\delta_{t-1}(i) a_{i j}\right] \\
\text { for } & 1 \leq j \leq N, 2 \leq t \leq T
\end{array}
$$

3) Termination:

$$
\begin{aligned}
& \bar{P}=\underset{1 \leq i \leq N}{\max } \quad \delta_{T}(i) \\
& \bar{q}_{T}=\begin{array}{l}
\arg \max \\
1 \leq i \leq N
\end{array} \quad \delta_{T}(i)
\end{aligned}
$$

4) Path (state sequence) backtracking:

$$
\bar{q}_{t}=\psi_{t+1}\left(\bar{q}_{t+1}\right) \quad t=T-1, T-2, \ldots, 1
$$

### 3.2.2 Modified Viterbi

Let us now relax the constraint that each segment corresponds to exactly one state. Because of the oversegmentation each state in a path through the HMM may now correspond to one, two, or three segments.

Let $\mathrm{s}_{\mathrm{ij}}=\left\langle\mathrm{s}_{\mathrm{i}}, \mathrm{s}_{\mathrm{i}+1}, \mathrm{~s}_{\mathrm{i}+2}, \ldots, \mathrm{~s}_{\mathrm{j}}\right\rangle$ be a word image subsequence. If a word image has R segments $\mathrm{s}_{1}, \mathrm{~s}_{2}, \ldots, \mathrm{~s}_{\mathrm{R}}$, the whole word can be written down as the image subsequence $\mathrm{s}_{1 \mathrm{R}}$.

The subsequence $s_{22}$ is just the sequence composed of only the second segment. We can regard each $\mathrm{s}_{\mathrm{ij}}$ as a subword that itself needs to be identified as a sequence of states. Denote the optimal state sequence for $\mathrm{s}_{\mathrm{ij}}$ as $\mathrm{Q}_{\mathrm{ij}}=\left\langle\mathrm{q}_{1}, \mathrm{q}_{2}, \ldots, \mathrm{q}_{\mathrm{T}}\right\rangle$, with probability $\mathrm{P}\left(\mathrm{Q}_{\mathrm{ij}}\right)$. $T \leq j-i+1$ in general, since word images are oversegmented. Recall that the states $q_{i} \in\{1,2$, ..., N \}, where N is the number of possible states.

Let $\alpha_{\mathrm{m}}\left(\mathrm{s}_{\mathrm{ij}}\right)$ be the optimal state sequence for $\mathrm{s}_{\mathrm{ij}}$ that begins with state m . Similarly, let $\omega_{\mathrm{n}}\left(\mathrm{s}_{\mathrm{ij}}\right)$ be the optimal state sequence for $\mathrm{s}_{\mathrm{ij}}$ that ends with state n , and let $\rho_{\mathrm{mn}}\left(\mathrm{s}_{\mathrm{ij}}\right)$ be the optimal state sequence that begins with $m$ and ends with $n$. The optimal state sequence $\mathrm{Q}_{\mathrm{ij}}$ is therefore $\rho_{m n}\left(s_{i j}\right)$ for some states $m$ and $n$.

Suppose we know the optimal state sequences $\omega_{\mathrm{n}}\left(\mathrm{s}_{1 \mathrm{k}}\right)$ and $\alpha_{\mathrm{m}}\left(\mathrm{s}_{\mathrm{k}+1, \mathrm{R}}\right)$ for all $1 \leq \mathrm{m} \leq$ $\mathrm{N}, 1 \leq \mathrm{n} \leq \mathrm{N}$, and $\mathrm{k}=1,2, \ldots, \mathrm{R}$. The optimal state sequence $\mathrm{Q}_{1 R}$ for the word $\mathrm{s}_{1 \mathrm{R}}$ is $\omega_{\bar{n}}\left(\mathrm{~s}_{1 \overline{\mathrm{k}}}\right) \alpha_{\overline{\mathrm{m}}}\left(\mathrm{s}_{\overline{\mathrm{k}}+1, \mathrm{R}}\right)$ where $\overline{\mathrm{k}}, \overline{\mathrm{n}}$, and $\overline{\mathrm{m}}$ are chosen by maximizing the probability

$$
P\left(Q_{1 R}\right)=\underset{1 \leq k \leq R, 1 \leq n \leq N, 1 \leq m \leq N}{\max } \quad P\left(\omega_{n}\left(s_{1 k}\right)\right) a_{n m} P\left(\alpha_{m}\left(s_{k+1, R}\right)\right)
$$

We recursively solve a similar optimization problem for $\mathbf{s}_{\mathbf{1} \overline{\mathbf{k}}}$ and for $\mathbf{s}_{\overline{\mathbf{k}}+\mathbf{1}, \mathbf{R}}$. In fact we can do this recursively for all $s_{i j}$. The recursion bottoms out when $j-i \leq 2$. If $i=j, P\left(Q_{i j}\right)=$ $\mathrm{P}\left(\mathrm{s}_{\mathrm{ij}} \mid \mathrm{c}\right)$ form some class c . That is, the optimal state sequences of the segment subsequence $\mathrm{s}_{\mathrm{ij}} \mathrm{is}$ just the neural network likelihood output. We have a similar case for $\mathrm{j}-1=1$ or $\mathrm{j}-1=2$, except that the new segment formed by the merged segments is just one of the possible paths.

To avoid duplicating work done during the recursion we do a bottom-up tabular approach. We first find $\rho_{m n}\left(s_{i j}\right)$ and $P\left(\rho_{m n}\left(s_{i j}\right)\right)$ for $\mathrm{j}-\mathrm{i}=0$, then for $\mathrm{j}-\mathrm{i}=1, \mathrm{j}-\mathrm{i}=2, \ldots, \mathrm{j}-\mathrm{i}=\mathrm{R}-1$. When $j-i=R-1$, we get $Q_{1 R}$ and we are done.

The algorithm is laid out more formally below. Remember that $\mathrm{P}\left(\rho_{\mathrm{mn}}\left(\mathrm{s}_{\mathrm{ij}}\right)\right)$ is the probability of a sequence and $\mathrm{P}\left(\mathrm{s}_{\mathrm{ij}} \mid \mathrm{c}\right)$ is a neural network output.

## Initialization:

We find the values of $\rho_{m n}\left(s_{i j}\right)$ and $P\left(\rho_{m n}\left(s_{i j}\right)\right.$ for chunks $s_{i j}$ where $j=i, j=i+1$, and $j=i+2$. These are chunks composed of one, two, or three contiguous segments, respectively. They are special because they can potentially represent single characters. Sequences of more than three contiguous segments can only be combinations of multiple characters.

The initialization, i.e. computation of $\rho_{m n}\left(s_{i j}\right)$ and $P\left(\rho_{m n}\left(s_{i j}\right)\right)$ for the smaller chunks, is more complicated than finding the same values for the larger subsequences because it involves neural network outputs.

For chunks with only one segment we have,

$$
\begin{aligned}
& \rho_{m, n}\left(s_{i, i}\right)=\left\{\begin{array}{cc}
\langle m\rangle & m=n \\
\text { nil } & \text { otherwise }
\end{array}\right. \\
& P\left(\rho_{m, n}\left(s_{i, i}\right)\right)=\left\{\begin{array}{cc}
P\left(s_{i, i} \mid m\right) & m=n \\
0 & \text { otherwise }
\end{array}\right. \\
& \text { for } \quad 1 \leq i \leq R, 1 \leq m \leq N, 1 \leq n \leq N
\end{aligned}
$$

For chunks with two segments we have,

$$
\begin{gathered}
\rho_{m, n}\left(s_{i, i+1}\right)= \begin{cases}\langle m\rangle & (m=n) \wedge\left(P\left(s_{i, i+1} \mid m\right)>P\left(s_{i i} \mid m\right) a_{m m} P\left(s_{i+1, i+1} \mid m\right)\right) \\
\langle m, m\rangle & (m=n) \wedge\left(P\left(s_{i, i+1} \mid m\right) \leq P\left(s_{i i} \mid m\right) a_{m m} P\left(s_{i+1, i+1} \mid m\right)\right) \\
\langle m, n\rangle & \text { otherwise }\end{cases} \\
P\left(\rho_{m, n}\left(s_{i, i+1}\right)\right)=\left\{\begin{array}{cc}
P\left(s_{i, i+1} \mid m\right) \quad \rho_{m, n}\left(s_{i, i+1}\right)=\langle m\rangle \\
P\left(s_{i i} \mid m\right) a_{m n} P\left(s_{i+1, i+1} \mid n\right) \quad \text { otherwise }
\end{array}\right. \\
\text { for } \\
1 \leq i \leq R, 1 \leq m \leq N, 1 \leq n \leq N
\end{gathered}
$$

For chunks with three segments, the possible $\rho_{\mathrm{mn}}\left(\mathrm{s}_{\mathrm{i}, \mathrm{i}+2}\right)$ are of the form $<\mathrm{m}>,<\mathrm{m}, \mathrm{n}>$, or $<\mathrm{m}, l, \mathrm{n}>$ where $1 \leq l \leq \mathrm{N}$. To simplify the expression for $\rho_{\mathbf{m n}}\left(\mathrm{s}_{\mathrm{i}, \mathrm{i}+2}\right)$, let us first define the probabilities of each of these candidates.

$$
\begin{gathered}
P\left(\langle m\rangle \mid s_{i, i+2}\right)=P\left(s_{i, i+2} \mid m\right) \\
P\left(\langle m, n\rangle \mid s_{i, i+2}\right)=\max \left(P\left(s_{i i} \mid m\right) a_{m n} P\left(s_{i+1, i+2} \mid n\right), P\left(s_{i, i+1} \mid m\right) a_{m n} P\left(s_{i+2, i+2} \mid n\right)\right) \\
P\left(\langle m, l, n\rangle \mid s_{i, i+2}\right)=P\left(s_{i i} \mid m\right) a_{m l} P\left(s_{i+1, i+1} \mid l\right) a_{l n} P\left(s_{i+2, i+2} \mid n\right)
\end{gathered}
$$

Now the optimal state sequence can be simply defined as

$$
\rho_{m, n}\left(s_{i, i+2}\right)=\left[\begin{array}{c}
\left.\rho=\begin{array}{c}
\arg \max \\
\langle m\rangle,\langle m, n\rangle,\langle m, l, n\rangle \\
1 \leq l, n, m \leq N
\end{array}\right] \quad\left\{P\left(\rho \mid s_{i, i+2}\right)\right\}
\end{array}\right.
$$

The probability of this path is therefore

$$
P\left(\rho_{m, n}\left(s_{i, i+2}\right)\right)=P\left(\rho_{m, n}\left(s_{i, i+2}\right) \mid s_{i, i+2}\right)
$$

## Recursive case:

$$
\begin{aligned}
\rho_{m, n}\left(s_{i, j}\right)= & {\left[\begin{array}{l}
i \leq k \leq j \\
1 \leq x \leq N \\
1 \leq y \leq N
\end{array}\right] \quad \rho_{m, x}\left(s_{i, k}\right) a_{x y} \rho_{y, n}\left(s_{k+1, j}\right) } \\
& \text { for } \quad 1 \leq m \leq N, 1 \leq n \leq N
\end{aligned}
$$

It is important to note that in order to avoid duplication of computations, we should do a bottom up approach where the current computation of $\rho_{\mathbf{m n}}\left(\mathrm{s}_{\mathrm{i}, \mathrm{j}}\right)$ depends only on values that have already been computed. First fill in all the values for the case where $\mathrm{j}-\mathrm{i}=3$ since all the $\rho$ 's for $\mathrm{j}-\mathrm{i}<3$ have been computed in the initialization. Then fill in the values for $\mathrm{j}-\mathrm{i}=$ 4, and so on.

## Termination:

When we have filled in the table for $\mathrm{j}-\mathrm{i}=\mathrm{R}-1$ from the recursive case and have found the $\rho_{\mathrm{mn}}\left(\mathrm{s}_{\mathrm{i}, \mathrm{R}}\right)$ paths, i.e. the optimal paths for the entire word $\mathrm{s}_{1 \mathrm{R}}$, it suffices to find the optimal choices of $m$ and $n$ to find the single best path,

$$
Q_{\text {opt }}=\underset{1 \leq n \leq N, 1 \leq m \leq N}{\max } \quad P\left(\rho_{m, n}\left(s_{1, R}\right)\right)
$$

### 3.2.3 Correctness of Modified Viterbi Algorithm

This section sketches a proof of correctness for the modified Viterbi Algorithm. We need to show partial and total correctness. The algorithm has partial correctness if, assuming it terminates, it comes up with the correct result. Total correctness is partial correctness combined with the fact that the algorithm terminates.

For partial correctness, we need to show that, if it terminates, the algorithm finds a state sequence that optimizes $\mathrm{P}\left(\rho_{\mathbf{m n}}\left(\mathrm{s}_{1, \mathrm{R}}\right)\right)$. To do this, we define the algorithm in terms of a generic dynamic programming algorithm and show that it has the optimal substructure property. This property arises from the fact that all the ways of splitting a subword $\mathrm{s}_{\mathrm{ij}}$ into two parts are themselves optimization problems. If they do not give optimal solutions, i.e. there exist solutions that have higher probability, then those solutions would be more optimal, thus posing a contradiction.

To show total correctness, it suffices to show that the algorithm terminates. The algorithm fills in the rows of the table, starting from row $r=1$ to row $r=R$. For each $r$ there are at most R-r table entries to fill in, and for each table entry, there are at most $\mathrm{N}^{2}$ combinations to consider. Since all these steps take a bounded number of computations, the entire algorithm itself take a bounded number of steps. Therefore, it terminates.

### 3.2.4 Asymptotic Performance

Let $T(r)$ be the time to find the paths $\rho_{m n}\left(\mathrm{~s}_{\mathrm{i}, \mathrm{i}+\mathrm{r}-1}\right)$ and their corresponding probabilities. Because we use a bottom-up approach, we can assume that the values for all subsequences of $\mathrm{s}_{\mathrm{i}, \mathrm{i}+\mathrm{r}-1}$ are already known.

$$
\left.\rho_{m, n}\left(s_{i, j}\right)=\begin{array}{c}
\arg \max \\
i \leq k \leq j \\
1 \leq x \leq N \\
1 \leq y \leq N
\end{array}\right] \quad \rho_{m, x}\left(s_{i, k}\right) a_{x y} \rho_{y, n}\left(s_{k+1, j}\right)
$$

There are $\mathrm{r}-1$ ways of choosing k in the optimization equation shown above. There are $N$ ways to choose $x$ and $N$ ways to choose $y$. Therefore, $T(r)$ is $O\left(r N^{2}\right)$.

The computation of the probabilities can be absorbed into this running time because we can store the values we find as we search for the argmax. In the base case where $\mathrm{r}=1$, the computation is also $\mathrm{O}\left(\mathrm{r} \mathrm{N}^{2}\right)=\mathrm{O}\left(\mathrm{N}^{2}\right)$ because this is the Neural Network output.

The algorithm goes through $r=1$, then $r=2$, and so on until $r=R$. There are $R$ subsequence sets in the row for $r=1, R-1$ subsequence sets for $r=2$, and only one for $r=R$. Thus, the number of subsequence sets for row $r$ is $R-r+1$. Hence the total runtime must be

$$
\begin{gathered}
\sum_{r=1}^{R}(R-r+1) O\left(r N^{2}\right) \\
=\left((R-1) \sum_{r=1}^{R} O\left(r N^{2}\right)-\sum_{r=1}^{R} O\left(r^{2} N^{2}\right)\right) \\
=\left((R-1) O\left(\left[\frac{R(R+1)}{2}\right] N^{2}\right)-O\left(\left[\frac{R(R+1)(2 R+1)}{6}\right] N^{2}\right)\right) \\
=O\left(N^{2}\left[\frac{(R+1)(R+1) R}{2}-\frac{R(R+1)(2 R+1)}{6}\right]\right) \\
=O\left(R^{3} N^{2}\right)
\end{gathered}
$$

In any particular instance of the algorithm, the number of states N is kept constant so the asymptotic running time is $\mathrm{O}\left(\mathrm{R}^{3}\right)$.

In practice, we speed up the algorithm by not keeping all $\mathrm{N}^{2}$ candidate subsequences. Many of the candidates have very low probabilities and are not worth keeping around.

Instead, only a short list is stored. This could be fixed in length or of variable length with a threshold for the probabilities for the elements on the list ${ }^{1}$.

### 3.3 Results.

Examples of input/output pairs of the modified Viterbi algorithm for HMMMOS are shown below. The hypothesis words are separated by vertical bars, the first being the top choice for the word and so on. Bear in mind that these are raw results which have not been used to generate candidates that actually belong to the lexicon. Furthermore, no context information has been used. Chapter 6 (results in section 6.5) introduces a scheme which uses context information and chooses hypotheses that can be found in the lexicon.

$$
\begin{aligned}
& 123 \text { Nath Glemroy } \\
& \text { Los Angeles, CA } 90049
\end{aligned}
$$

\{ $123|\mathrm{a}| \mathrm{w}$ \} \{ wenglenwy | naaglenwy | wenghiny \}
$\{60 \mid$ bo $\mid$ la $\}\{$ age, $\mid$ ashe, $\mid$ ass 7$\}\{$ ca $\mid$ la $\mid$ ea $\}\{9 \mathrm{ma} \mid$ ama $\mid$ ana $\}$
Analysis: The first word, the number " 123 ", is correctly recognized as the first hypothesis. The words "North Glenroy" are combined into a single word by the algorithm because of their proximity. The "glen" part is captured by all three hypotheses, as is the " $y$ " at the end. The "ro" is always recognized as a " $w$ ", and this is consistent with a human inspection of the letters as drawn. "North" is not satisfactorily recognized although it is hoped that given the context of the other letters, the words "North Glenroy" becomes one of they candidate words. In the second line, "los" is recognized as " 60 ", "bO", or "la"

[^1]because of the shape. The hypothesis "agle" captures four of the letters in "angeles" correctly and "CA" is correctly recognized. Due to the cursive manner in which the zip code is written, it is recognized as a sequence of alphabetic characters instead of numeric digits.

$\{345|34 \mathrm{~s}| \mathrm{b} 45$ \} ~ \ { a | w | a n ~ \ } ~ \ { ~ a n e ~ | ~ p i n e ~ | ~ a v e ~ \ } ~ \ { s t u ~ | ~ m n ~ | ~ s t r ~ \ } ~
\{ wang | wang | awiling \} ~ \ { ~ w a n ~ | ~ s a n a ~ | ~ s a n ~ \ } ~ \ { ~ o 3 9 3 7 ~ | ~ o 3 9 5 7 ~ | ~ o 3 9 5 9 ~ \ } ~
Analysis: The first word " 345 " is accurately recognized. The second word, which the author cannot read himself, is hypothesized to be " $a$ ", " $w$ ", or "an". The correct reading of the next word, "pine", is actually the second hypothesis for that word. "st" is interpreted as "stu" because in this instance of the HMMMOS "." is not a character. In the second line, "Reedsburg" is not satisfactorily recognized, and the author cannot read the state. The last zip code hypothesis has the correct digits except for the leading " 5 ".
\[

$$
\begin{aligned}
& 1000 \text { N. Market st. } \\
& \text { Milwaukee, }{ }_{5 I} 53202
\end{aligned}
$$
\]

$\{1000|100|, \mathrm{m}\}\{\mathrm{n} 7|\mathrm{nn}| \mathrm{w} 7\}\{$ nundsta $\mid$ nalatsta $\mid$ nundsto $\}$
\{nllwadia, | nllwadale, | mllwadia, \} ~ \ { ~ w i ~ | l v i ~ | ~ w t ~ \ } ~
$\{5|\mathrm{~s}| 3\}\{32 \mathrm{o} 2|\mathrm{ao} 2| 22 \mathrm{o} 2\}$

$$
\begin{aligned}
& \text { P.O. BOX } 354 \\
& \text { WAUPACA, WI } 54981 \\
& \{\mathrm{pa}|\mathrm{ra}| 8 \mathrm{o}\}\{\text { box } \mid \text { bor } \mid 8 \mathrm{ox}\}\{35|3 \mathrm{~s}| 3 \mathrm{a}\}\{4|\mathrm{t}| 7\} \\
& \text { \{ wau | wav \| wal \} \{ paca, | nca, | raca, \} \{ vil wilui \} \{ } 549|54 \mathrm{~g}| 5 \mathrm{t} 9\}\{81|\mathrm{al}| \mathrm{pl}\} \\
& 7457 \text { blmurnel Aue. } \\
& \text { Aveideletow. W1 53562 } \\
& \text { \{ } 74|14| \text { ta }\}\{57|a| w\}\{\text { ewinanaven | ewintwaven } \mid \text { ewinawaven \} } \\
& \{\text { misileton | wisileton | mosileton \} \{, |1|i\} \{wi|w, lo51\}\{5302|5330|5350 \} } \\
& 53716 \\
& \{5 \mid \text { brl } 52\}\{40|50| \text { ho }\}\{y|9| g\}\{\text { may }|m y| n e y\}\{1004|m n t,| a n o 4\}\{4|a| a l\} \\
& \{\text { mononn } \mid \text { monoun } \mid \text { monann }\}\{,|1| r\}\{\text { wilwl|win }\}\{53716|o 7716| a y 6\} \\
& \text { \{pna|rna|rda\} \{428|423o|pa\} }
\end{aligned}
$$

$\{$ bla $\mid$ blan $\mid$ ma $\}\{$ acad, $\mid$ mad, $\mid$ wad, $\}\{n|h| w\}\{n y|c y| w y\}$
\{ 11545|115457|11575 \}

$$
\begin{aligned}
& \text { PO.Box } 568 \\
& \text { Honesdele, PA } 18431
\end{aligned}
$$

\{80xn $|8080 \mathrm{x}|$ poxn \} $\{5 \mathrm{kf}|5 \mathrm{hf}|$ shf $\}$
\{ nannd, | nannn, I nanwa, \} \{ pa|ra|ta \} $16431|18431|$ inn \}

$$
\begin{aligned}
& \text { 9in }{ }^{\text {Jillizusson St. }} \\
& \text { Madison, WI 53703-3549 }
\end{aligned}
$$

\{ 917 |947|91n \} \{ willimon | willianon | willianson \} \{ st | st, | sti \}
\{ madgonlwi | madigonlwi | madgonlwl \} \{ 537o3-3549 | 5303-3549 | mo3-3549 \}

\{3|37|57\}\{15|li|i5\}\{n|w|v\}\{heng|heno|heny\}\{st|a|w\}
\{ nadson, | madson, | nadsoni \} \{ wi | we | all \} \{ $53703|5323| 53233$ \}

## Chapter 4

## Probability Estimation

### 4.1 Overview

Two kinds of probability distributions are used in the Hidden Markov Model. There are transition probabilities for the underlying Markov process and there are observation probabilities that map chunks of word images to the "hidden" states. The estimation of both types of distributions are discussed in the sections that follow.

### 4.2 Initial and Transition Probabilities

The probability that any word will start at a particular character or state $i$ is called the initial probability $\mathrm{a}_{\pi, \mathrm{i}}$. A transition probability $\mathrm{a}_{\mathrm{ij}}$ is the probability of going from state i to state j in the word sequence. I.e., the probability that character i is immediately to the left of character j in any word sequence. Initial and transition probabilities can be estimated from a lexicon. In the absence of a separate lexicon, the training data could be used.

### 4.3 Observation Likelihood Estimation (Character recognition)

We try three techniques for estimating the or observation probability distributions: a Multilayer Perceptron (MLP), a recurrent network, and a Hierarchical Mixture of Experts (HME).

Neural Networks approximate a posteriori probabilities $\mathrm{P}(\mathrm{jls})$. However, we need a neural network to estimate likelihoods $\mathrm{P}(\mathrm{s} \mid \mathrm{j})$. A conversion from a posteriori probabilities to likelihoods is given by Bayes rule. If $s$ is the segment image and $j$ is the character class being considered, then

$$
P(s \mid j)=\frac{P(j \mid s) P(s)}{P(j)}
$$

For a particular segment s , the distribution of likelihoods for all j is proportional to the following expression.

$$
P(s \mid j) \approx \frac{P(j \mid s)}{P(j)}
$$

The relative magnitudes of these likelihoods are preserved in the approximation above. The actual likelihoods can therefore be computed by normalizing the estimates. Note that $\mathrm{P}(\mathrm{j})$ can be readily estimated from the training data.

### 4.3.1 Multilayer Perceptron (MLP)

The multilayer perceptron, trained using backpropagation, is a popular neural network architecture for this particular application. The appeal of this technique is in its simplicity in implementation and the fact that it has been successfully applied to various handwriting recognition problems in the past.


Figure 4.1: Multilayer Perceptron

MLPs consist of simple calculation elements, called neurons, and weighted connections between them. In a feedforward Multilayer Perceptron (MLP) the neurons are
arranged in layers. A neuron from one layer is fully connected with the neurons of the next layer.

The first and last layers are the input and output layers, respectively. The layers between them are called hidden layers. Feature vectors are given to the neurons in the input layer; the results come out of the output layer. The outputs of the input neurons are propagated through the hidden layers of the net. The figure below shows graphically the algorithm that each neuron performs.


Figure 4.2: Node at layer $h$ and row $j$

The activation $\mathrm{z}_{\mathrm{hj}}$ of a hidden or output neuron j at layer h is the sum of the incoming data multiplied by the connection weights like in a matrix product. The individual bias value $r_{h j}$ is then added. I.e.

$$
z_{h j}=r_{h j}+o_{h-1,1} w_{h-1,1}+o_{h-1,2} w_{h-1,2}+\ldots+o_{h-1, k} w_{h-1, k}
$$

There are k neurons in layer $\mathrm{h}-1$. The output $\mathrm{u}_{\mathrm{hj}}$ is calculated by a nonlinear function f. A popular nonlinearity used in practice is the sigmoid function.

$$
u_{h j}=f\left(z_{h j}\right)=\frac{1}{1+e^{-z_{n j}}}
$$

A feedforward multilayer perceptron can approximate any function after a suitable
amount of training. Known discrete values of this function, together with expected outputs, are presented to the net in the training phase. It is then expected to learn the function rule [Rumelhart et. al. 1994].

The behavior of the net is changed by modification of the weights and bias values. The network is trained using the Backpropagation algorithm [Cun 1986] [Rumelhart et. al. 1986].

A common performance metric for neural network training is the sum of squared errors. Although theoretically the outputs of a MLP trained using the sum of squared errors can approximate a posteriori probabilities [Richard \& Lippmann 1991] (under certain assumptions such as a sufficient number of training samples), in practice, due to the finite number of training samples, the probability distributions for likelihoods of character classes given the character image are quite sharp. That is, there are one or two character classes with very high probabilities and the other character classes have close to zero probabilities. Due to high variability in handwriting, the true character class may be given a likelihood that is very close to zero. We want a smoother distribution that will make it more likely for a globally oriented dynamic programming algorithm to pick the true character class in spite of it having a low likelihood for a particular instance. We propose to compare the results between using the Kullback-Leibler divergence and the sum of squared errors. The paper [A2iA 1996] discusses how probability distributions can be smoothened using the combination of the softmax function and the Kullback-Leibler divergence as a metric.

### 4.3.2 Recurrent Neural Net

A Recurrent Network was also tried. The network takes in two feature vectors: that of the chunk being recognized as a character and that of the segment immediately to the left. The latter provides context which would hopefully help the network identify the chunk correctly. The context helps distinguish between character classes which look similar by themselves but look obviously different when they are seen adjacent to the preceding segment.


Figure 4.3: Right and middle parts of " $w$ " vs. " $u$ " when context is present.

For example, the chunk composed of the middle and right ascenders of a " $w$ " looks very similar to a whole "u" (See Figure 4.3). We would like the chunk to be negatively recognized as a " $u$ " and for a real " $u$ " to be recognized as a " $u$ ". If we include the preceding segment of each, we get a " $w$ " and a "au", respectively, where " $a$ " is some other character. Now it is very clear that the two images are different.

The basic architecture of the recurrent net that was implemented is a Multilayer Perceptron. The network operates exactly as an MLP except that the feature vector for an extra segment is used.

In order to provide even more context, the segment immediately to the right was also
included. The resulting architecture is a double-sided network.

### 4.3.3 Hierarchical Mixture of Experts

The paradigm of divide-and-conquer is a prominent one in computer science. By dividing a problem into smaller subproblems and then combining the results, we can often come up with efficient and elegant solutions. Divide-and-conquer can be applied to the neural network approach to handwriting recognition by dividing the feature vector space and having a neural network specialize on each division. Other neural networks can be used to arbitrate between these specialists. The combinations of the arbitrators and the specialist neural networks are called modular networks. An interest in modular networks has been stimulated in recent years by the emergence of 'space partitioning algorithms' like Classification and Regression Trees (CART) [Breiman, Friedman, Olshen \& Stone 1984], Multivariate Adaptive Regression Splines (MARS) [Friedman 1991], and Hierarchical Mixtures of Experts (HME) [Jordan \& Jacobs 1994].

The divide-and-conquer approach embodied by modular neural networks is an intriguing one for handwriting recognition. People have various styles in handwriting but there are certain regularities in handwriting styles within certain groups. For example, children tend to draw bigger and more irregular characters than adults. The writer's nationality and sex may affect the character shape, stroke width, and slant in somewhat predictable ways. In light of these observations, we try to use a modular neural network for off-line cursive handwriting recognition. In particular, we implemented the Hierarchical Mixture of Experts (HME) model.

Among the modular neural networks mentioned above, HME is particularly attractive because it uses "soft splits" of the data. This means the data is allowed to lie simultaneously in multiple regions. This reduces the variance of the recognizer and therefore helps correct one of the problems associated with dividing the input space. The HME has already been applied to regression and classification problems such as robotics control [Jordan \& Jacobs 1994] and speech recognition [Waterhouse 1993]. It has not, however, been applied to Handwriting Recognition.

The evolution of HMEs has been enhanced by the use of Generalized Linear Models (GLIMs) in the network and the development of a fast algorithm for training the 'Hierarchical Mixture of Experts' using the Expectation Maximization (EM) Algorithm of [Dempster, Laird, \& Rubin 1997] [Jordan \& Jacobs 1994]. The EM algorithm converges significantly faster than the more familiar gradient descent algorithms and could improve the convergence times for our particular application.

The HME architecture is shown in Figure 4.4. It is a tree with gating networks at the nonterminals and expert networks at the leaves. Gating networks receive each feature vector x as input and produce scalar outputs that are a partition of unity. Each expert network outputs a vector $\mu_{\mathrm{ij}}$ for each input feature vector x . At the bottom-most level of nonterminals, the $\mu_{\mathrm{ij}}$ are scaled by the values of the outputs of their corresponding gating networks to form intermediate outputs for the next higher level of nonterminals. This process continues until the root of the tree is reached. The output $\mu$ at the root is the output of the entire HME tree.


Figure 4.4: A 2-level HME. To form a deeper tree, each expert is expanded recursively into a gating network, a nonterminal node, and a set of subexperts.

### 4.4 Results

The character-wise performance for the MLP is

Training patterns: 158898 characters
Accuracy: 77.50\% (top choice)
Test patterns: 25300 characters

Accuracy: $74.30 \%$ (top choice)
The architecture of the MLP is

108 input units corresponding to 108 features.

27 outputs for the Alpha network (26 letters and a reject class).
80 hidden units are used.
The HME never converged when Generalized Linear Models (GLIMs) were used because the covariance matrices became singular very rapidly and the training algorithm
diverged. Attempts to control the divergence by enforcing absolute lower bounds on the eigenvalues did not help.

The recurrent network (along with the one with pre- and post-context) also did not work properly.

## Chapter 5

## Rejection of bad samples to increase recognition rates

### 5.1 Problem Definition

The system described in previous chapters does not correctly recognize all inputs. There are inputs that are too noisy, too distorted, or just do not lend themselves to proper preprocessing. In order to increase the recognition rate we can reject these bad inputs.

The situation can be modeled by having two distributions of oversegmented word images. One is for all word images that will be correctly recognized by our system. The other is for word images that will not. The distributions live in a multidimensional feature space ${ }^{1}$. Our goal is to increase the recognition rate by identifying a decision boundary and using our recognition system only on those that are on the "good" side of this boundary.

A simple example of this model is when there are two one-variable Gaussian distributions in a random variable x as shown in Figure 5.1.

Distribution A is for the bad inputs; B is for the good inputs, i.e. the ones that will be correctly identified. $\mathrm{p}_{1}$ is the probability of an input being bad; $\mathrm{p}_{2}$ is the probability of an input being good. $\mathrm{p}_{1}+\mathrm{p}_{2}=1$. Since our space consists of only one scalar variable x , the decision boundary consists of a single threshold $\mathrm{x}_{0}$. In our example, every word image whose feature x is greater or equal to $\mathrm{x}_{0}$ is used as input into the system for recognition. Each image whose feature x is less than $\mathrm{x}_{0}$ is rejected.

[^2]

Figure 5.1: Two gaussians representing good and bad word images.

Some bad images might still be included in the input set for the system because their feature values are on the "good side" of the boundary. The probability of including a bad image in the input set is $\alpha$ as shown in Figure 5.1. It is usually called the false positive rate. Similarly, some good images will be rejected, causing the recognition rate of the system to go down. The probability of rejecting a good image is $\beta$ as shown. It is called the miss rate. The rejection rate of the system is defined as $p_{1}(1-\alpha)+p_{2} \beta$. The recognition rate is $p_{2}(1-\beta) /\left[p_{1} \alpha+p_{2}(1-\beta)\right]$. The overall success rate is $p_{2}(1-\beta)$.

It is desirable to have a high recognition rate and a low rejection rate. $x_{0}$ can be moved around to manipulate these values but there is a trade-off. In most cases a lower $\alpha$ increases the recognition rate of the system. As we move $x_{0}$ to the right, $\alpha$ decreases but $\beta$
increases, thus increasing the rejection rate. The opposite happens when we move $\mathrm{x}_{0}$ to the left. $\alpha$ increases causing the recognition rate to decrease. However the rejection rate decreases at the same time due to the decreased $\beta$. To maximize the overall success rate, the optimal value of $x_{0}$ has to be found. The next section discusses our method of finding this value.

### 5.2 Multilayer Perceptron Approach

In practice the distributions are not Gaussian and x is a multi-dimensional vector, with $\mathrm{x}_{0}$ representing a surface in that space. Finding the optimal $\mathrm{x}_{0}$ is equivalent to finding the best partitioning of the space into two parts: the good and the bad side. We do this using a Multilayer Perceptron (MLP).

An input to the rejection system is composed of an oversegmented word image and its corresponding candidate word list. The verification algorithm discussed in Chapter 6 is performed to rank the candidates. Features are then extracted from all this information and fed into the rejection MLP. The output is a boolean value which is normally represented by real number with a threshold. If the number is greater than or equal to zero, the word image is considered to be on the good side and is kept as input. If it is less than zero, it is pronounced bad and is rejected.

### 5.3 Features used

The nine features that were used as input to the MLP are listed below. When it is not specified, the feature refers to the top candidate.

1. Probability of top candidate divided by the number of characters of the top candidate.
2. Difference in probability of the top two candidates.
3. Number of deleted segments divided by the number of characters.
4. Number of characters that were ranked no. 1 by the neural character classifier divided by the number of characters.
5. Number of characters that were ranked no. 2 by the neural character classifier
divided by the number of characters.
6. Number of characters that were ranked no. 3 by the neural character classifier divided by the number of characters.
7. Number of characters that were ranked no. 4 by the neural character classifier divided by the number of characters.
8. Number of characters that were ranked no. 5 by the neural character classifier divided by the number of characters.
9. Number of characters that were ranked no. 6 to 10 by the neural character classifier divided by the number of characters.

Feature 1 distinguishes hypotheses that have high probability purely because there are fewer characters, and hence fewer neural net likelihoods to multiply, from those that have high probabilities because the average neural net likelihood is high. This feature is intended to salvage very long hypothesis words.

From empirical results, the difference in probability between the true hypothesis and any other hypothesis is large compared to the difference between two wrong hypotheses. Feature 2 arises from this observation.

Feature 3 separates hypotheses that have many delete characters from those with few. A delete character is a one-segment character that is unreadable and does not belong to a chunk that is recognized as a character, i.e. it does not represent any of the states of the HMM. A good match between a hypothesis and the word image should not create too many delete characters relative to the number of characters in the hypothesis.

Features 4 to 9 collect information on the relative rankings of the characters in the hypothesis based on neural network likelihoods. The higher the rankings of the character states in the hypothesis, the more confident we are that the hypothesis is a good match for the word image.

### 5.4 Results

Figure 5.2 shows the trade-off between the error rate and the rejection rate when a rejec-
tion network is used to filter word images that are fed into the verification algorithm. The data is taken from the CEDAR (Center for Excellence for Document Analysis and Recognition) database. All the usable city name images were used.


Figure 5.2: Error rate v.s Rejection rate for the verification algorithm

## Chapter 6

## Improving recognition by verification and correction

### 6.1 Introduction

The system outlined in Chapter 3 produces sequences of characters which might not even correspond to actual words. This presents two opportunities to improve the recognition rate: (a) the lexicon can be used to identify candidate sequences that are also actual words; and (b) the context given by surrounding word images could help better identify the best state sequence for each word image.

This chapter discusses a new dynamic programming algorithm which exploits these opportunities. It ranks word candidates taken from the lexicon for a particular segmented word image. This constrains the list of candidates to be members of the lexicon. The candidates are chosen using an existing verification and correction system [Lorie 1994] that finds words (state sequences) that are actually in the lexicon and "fit" the neighboring words.

The input to the verification and correction system is a list of candidate state sequences for each word image in an entire line of text. This is normally generated by the system described in Chapter 3. The output is a list of candidate state sequences for each word image. These candidates are from the lexicon and were chosen for their closeness to the input state sequences and for their relevance in the context provided by the neighboring word images.

Once the candidate list is retrieved, the candidates are ranked by the dynamic programming algorithm. The algorithm, using the oversegmented word image, maps the candidates to probabilities and sorts them according to the probabilities.

### 6.2 A Dynamic Programming Algorithm

In this section we describe a new dynamic programming algorithm that, given an oversegmented word image, maps candidate state sequences to probabilities.

Let us review some notation. A word image is of the form $S=\left\langle s_{1}, s_{2}, \ldots, s_{R}\right\rangle$. The $s_{i}$ are segments and there are R of them. A candidate is of the form $\mathrm{Q}=<\pi, \mathrm{q}_{1}, \mathrm{q}_{2}, \ldots, \mathrm{q}_{\mathrm{T}}>$. There are T character states. Because the image is oversegmented, each character could be represented by at least one and at most three contiguous segments. Therefore $T \leq R$.

We introduce a new construct, combo, which is a partition of S into chunks. Recall that a chunk is a sequence of one, two, or three segments contiguous segments. A combo is of the form $\mathrm{O}=<\mathrm{o}_{1}, \mathrm{o}_{2}, \ldots, \mathrm{o}_{\mathrm{T}}>$ and it forms a one-to-one mapping between chunks and non$\pi$ states in the candidate. The chunk $\mathrm{o}_{1}$ can only be one of $\left\langle\mathrm{s}_{1}\right\rangle,\left\langle\mathrm{s}_{1}, \mathrm{~s}_{2}\right\rangle$, and $\left\langle\mathrm{s}_{1}, \mathrm{~s}_{2}, \mathrm{~s}_{3}\right\rangle$. Similarly, $\mathrm{o}_{\mathrm{T}}$ could be $\left\langle\mathrm{s}_{\mathrm{T}-2}, \mathrm{~s}_{\mathrm{T}-1}, \mathrm{~s}_{\mathrm{T}}\right\rangle,\left\langle\mathrm{s}_{\mathrm{T}-1}, \mathrm{~s}_{\mathrm{T}}\right\rangle$, or $\left\langle\mathrm{s}_{\mathrm{T}}\right\rangle$. The combo is only valid if, (a) each segment occurs in at most one chunk; (b) the combo contains all the segments in the word image; and (c) if a segment appears before another segment in the word image, the chunk of the former appears before, or in the same chunk as, the latter.

If $T<R$, multiple combos correspond to the word image. Let us define the function $f$ over the set of all oversegmented words $S$ to be $f(S)=\{O \mid O$ is a valid combo for $S\}$. If $S$ corresponds to the word $Q$, all but one of the elements $O$ of $f(S)$ are mis-segmentations of Q. Otherwise, all the elements of $f(S)$ will be mis-segmentations. At least one chunk $o_{i}$ of a mis-segmentation combo will not look like a valid character. A few elements of $f(S)$ will prove to be a close match for the candidate. The degree to which a combo is close to the perfect segmentation is captured in a probability estimate. If the state sequence $Q$ is a bad candidate for S , all the probabilities for the O 's will be low. If Q is a good candidate, at least one candidate will have a high probability.

The algorithm discussed here finds the combo with the highest probability. This probability, multiplied by the probability of Q itself, is then regarded as the probability mapped to the candidate in the system described in section 6.1.

The a posteriori probability estimate for a combo O given a candidate Q is

$$
P(O \mid Q)=P\left(o_{1} \mid q_{1}\right) P\left(o_{2} \mid q_{2}\right) \ldots P\left(o_{T} \mid q_{T}\right)
$$

The best combo $\mathrm{O}^{*}$ and its probability are given by

$$
\begin{gathered}
O^{*}=\begin{array}{c}
\arg \max \\
O \in f(S)
\end{array} \quad P(O \mid Q) \\
P(O, Q)=P\left(O^{*} \mid Q\right) P(Q)
\end{gathered}
$$

We have taken out all the conditioning with respect to $S$ in the expressions above since the equations are all conditioned on some word image $S$. Recall that $P(Q)$ is just the Markov chain probability

$$
P(Q)=a_{\pi, q_{1}} a_{q_{1}, q_{2}} \ldots a_{q_{r_{-1}}, q_{T}}
$$

So now the problem is how to find $\mathrm{O}^{*}$ efficiently. We go through the algorithm with an example. Suppose we had the word image in Figure 6.1 and we wanted to match it with the state sequence $\mathrm{Q}=\langle\pi$, ' $w$ ', ' $o$ ', ' $r$ ', ' $d$ ' $>$.


Figure 6.1: An oversegmented image of "word".

The dynamic programming table for this case is shown in Figure 6.2.


Figure 6.2: Dynamic Programming Table $\psi$ for "word".

Every combo must start at the $(\pi, \pi)$ position and end at the lower right square on the table. Each chunk corresponds to an arrow. An arrow goes from one row to the next, and down and to the right. It can only go one, two, or three spaces to the right, and one space down. The number of spaces it goes to the right denotes the number of segments the chunk eats up.

For example, in the combo shown in Figure 6.2, the chunk $\mathrm{o}_{1}=\left\langle\mathrm{s}_{1}, \mathrm{~s}_{2}, \mathrm{~s}_{3}\right\rangle$ corresponds to ' $\mathrm{w}^{\prime}$ ', $\mathrm{o}_{2}=\left\langle\mathrm{s}_{4}\right\rangle$ corresponds to ' $\mathrm{o}^{\prime}$ ', $\mathrm{o}_{3}=\left\langle\mathrm{s}_{5}, \mathrm{~s}_{6}>\right.$ corresponds to ' r ', and $\mathrm{o}_{4}=\left\langle\mathrm{s}_{7}\right\rangle$ correspond to ' $d$ '. The sequence $<\pi, o_{1}, o_{2}, o_{3}, o_{4}>$ happens to be the optimal combo. The table is used to explore different possible combos and find the optimal (highest probability) one.

Note that the algorithm does not use the entire table. There are entries that are useless because they violate the constraints on the relationship between chunks and segments. Therefore, only squares within the parallelogram are used.

Let us now go through the algorithm in detail. It requires two dynamic programming tables: $\operatorname{LP}(i, j)$ and $\psi(i, j)$, where $i=0,1, \ldots, T$ and $j=0,1, \ldots, R$. Each element in the table
corresponds to a state i and a segment j . As a special case $\operatorname{LP}(0,0)$ and $\psi(0,0)$, the upper left square in the tables, correspond to the $\pi$ vs. $\pi$ element. An example of a $\psi$ table is shown on Figure 6.2.

For each valid $i$ and $j$ value, $\operatorname{LP}(i, j)$ is to the maximum $\log$ probability ${ }^{1}$ of ending the $\mathrm{i}^{\text {th }}$ character with the $\mathrm{j}^{\text {th }}$ segment, and hence the acronym. Each element of the table $\psi(\mathrm{i}, \mathrm{j})$ is a pair $<\mathrm{k}, \mathrm{l}>$ which indicates the position of the box which maximizes the probability LP(i,j).

## Initialization:

1. For $i=0,1, \ldots, T$
2. $\operatorname{For} \mathrm{j}=\mathrm{i}, \mathrm{i}+1, \ldots, \min (\mathrm{i}+\mathrm{T}-1, \mathrm{~T})$
3. $\quad L P(i, j)=-\infty$
4. $\quad \psi(\mathrm{i}, \mathrm{j})=$ nil

The initialization phase sets all the interesting elements of LP to negative infinity, the lowest $\log$ probability possible, and all the interesting elements of $\psi$ to nil.

Fill in the tables:
5. $\operatorname{LP}(0,0)=0$
6. For $\mathrm{i}=1, \ldots, \mathrm{~T}$
7. $\operatorname{For} \mathrm{j}=\mathrm{i}, \mathrm{i}+1, \ldots, \min (\mathrm{i}+\mathrm{T}-1, \mathrm{~T})$
8. $\quad$ For $k=\max (i, j-3), \max (i, j-3)+1, \ldots, j-1$
9. $\quad \operatorname{tmp}=\operatorname{LP}(\mathrm{i}-1, \mathrm{k})+\log (\mathrm{NN}(\mathrm{k}+1, \ldots, \mathrm{j}))$
10. If $(\operatorname{tmp}>\operatorname{LP}(\mathrm{i}, \mathrm{j}))$
11. $\quad L P(i, j)=t m p$
12. $\quad \psi(\mathrm{i}, \mathrm{j})=<\mathrm{i}-1, \mathrm{k}\rangle$

1. Note that maximizing log probabilities is the same as maximizing probabilities since logarithm is a monotonically increasing function.

Line 5 sets the probability of $\operatorname{LP}(0,0)$ to zero in order to give it the highest $\log$ probability among other squares on the 0 'th row, making it the starting point of all paths through the table. The next row will only point to either $\langle 0,0>$ or to nil. Line 6 starts the loop that goes through the rows of each table going from the first to the last $\left(\mathrm{T}^{\text {th }}\right)$ character in the hypothesis word. Line 7 starts the loop through the columns in each row. In order to have only valid combos as discussed earlier, not all the columns are explored. The loop on Line 8 goes through the possible ways a link could point to the box being considered. There are at most three, each corresponding to the number of segments that are joined together to form a chunk. For each of these possibilities, the log probability of the source box is added to the neural net output for the corresponding chunk (see line 9). The neural net takes in the feature vector extracted from the chunk composed of the $\mathrm{k}+1^{\text {st }}$ up to the $\mathrm{j}^{\text {th }}$ segment. The source box that gives the highest $\log$ probability in line 9 is chosen and the $\log$ probability value is updated in line 11 . Line 12 stores the location of this box in the $\psi$ table.

The optimal log probability for the hypothesis is $\operatorname{LP}(T, R)$. We often want to find the optimal combo that gave rise to this probability to visually check the algorithm's behavior. This is done by tracing back the pointers in the $\psi$ table as described in the pseudocode below. The notation $\psi(\mathrm{x}, \mathrm{y})$.j gives the j -component of the tuple stored in $\psi(\mathrm{x}, \mathrm{y})$.

$$
\text { 13. } \mathrm{C}_{\mathrm{T}}=\psi(\mathrm{T}, \mathrm{R}) \cdot \mathrm{j}
$$

14. For $i=T-1, T-2, \ldots, 1$
15. $\quad C_{i}=\psi\left(i, C_{i+1}\right) . j$

Each $\mathrm{C}_{\mathrm{i}}$ is one less than the index of the first segment of the $\mathrm{i}^{\text {th }}$ chunk. It could also be interpreted as the index of the last segment of the $i-1{ }^{\text {th }}$ chunk. In line $13, \mathrm{C}_{\mathrm{T}}$ is set to be the j-component of the $\psi(T, R)$, i.e. the segment or column number of the square that maximizes the step to the last chunk. It is therefore the index of the last segment of the chunk that precedes the last chunk in the sequence. In lines 14 and 15 , the processes recurses by
finding the j -component of the square that maximizes the step to the last chunk on the reverse path. $\mathrm{C}_{1}$ must be equal to zero because all optimal paths start at the upper left hand square.

The optimal combo is therefore

$$
\left\langle\left\langle s_{C_{1}+1}, \ldots, s_{C_{2}}\right\rangle,\left\langle s_{C_{2}+1}, \ldots, s_{C_{3}}\right\rangle, \ldots,\left\langle s_{C_{T}+1}, \ldots, s_{T}\right\rangle\right\rangle
$$

The algorithm does not allow for the existence of delete characters, segments that are specks of noise in the bitmap and do not really correspond to a valid character. Our implementation extends the algorithm above by allows for one-segment delete characters by eating up a segment without assigning it to a character/state. This implies a horizontal movement in the dynamic programming tables. In the pseudocode, line 8 is extended to also go through the box directly to the left of the box being considered, i.e. $<\mathrm{i}, \mathrm{j}-1>$. The neural network probability assigned to this single segment is its likelihood of being a delete character divided by an empirically derived parameter. This parameter is introduced because most single segments would have a high probability of being a delete character and delete characters should only be used sparingly to avoid mistaking parts of actual characters for delete characters. Note that the extra processing introduced in lines 8 to 12 adds a greedy component in the algorithm because local decisions are made about the identity of a segment, i.e. being a delete character.

Another exception to the basic algorithm is the case where the oversegmentation fails and two characters are actually put together in a single segment. This happens rarely but if not handled properly, it causes the algorithm to perform suboptimally. Fixing it involves considering box $<\mathrm{i}-2, \mathrm{j}-1>$ in addition to the other boxes in line 8 . This modification does not introduce any greediness to the algorithm.

Note that these two modifications expand the parallel work area into the entire rectangular table for both LP and $\psi$. This implies an increase in the running time.

### 6.3 Correctness

For the purpose of proving the correctness of the dynamic programming algorithm consider only the basic algorithm which does not have any greedy components.

For partial correctness, it suffices to show that the algorithm has an optimal substructure property. At each box, the values that are compared must be optimal themselves. Otherwise, there is a contradiction.

The algorithm terminates because there are a finite number of entries in the tables to fill. Hence, we have total correctness.

### 6.4 Asymptotic Performance

We do an $\mathrm{O}(1)$ computation for each element of the parallelogram in the table. The width of the parallelogram is $\mathrm{R}-\mathrm{T}+1$. The length is $\mathrm{T}+1$. So the parallelogram has $(\mathrm{R}-\mathrm{T}+1)(\mathrm{T}+1)$ elements and filling up the table takes $\mathrm{O}(\mathrm{R}-\mathrm{T}+1)(\mathrm{T}+1) \mathrm{O}(1)$ or just $\mathrm{O}\left(\mathrm{RT}-\mathrm{T}^{2}\right)$. Backtracking takes $\mathrm{O}(\mathrm{R}-\mathrm{T}+1+\mathrm{T}+1)=\mathrm{O}(\mathrm{R})$ time so filling up the table so the whole algorithm takes $\mathrm{O}\left(\mathrm{RT}-\mathrm{T}^{2}\right)+\mathrm{O}(\mathrm{R})$. Informally, since the backtracking path contains unique blocks and these blocks form a proper subset of the parallelogram, the total running time must be O (RT-$\left.T^{2}+\mathrm{R}\right)=\mathrm{O}\left(\mathrm{RT}-\mathrm{T}^{2}\right)$. Recall that $\mathrm{T} \leq \mathrm{R}$ because there are at most as mary states as there are segments. Therefore, an upper bound on the running time is $O\left(R^{2}\right)$.

If we use the entire table, as is the case when we allow for delete characters and segments which contain more than one character, the running time is $\mathrm{O}(\mathrm{RT})$. This is bounded above by $\mathrm{O}\left(\mathrm{R}^{2}\right)$.

### 6.5 Results

The rejection network was evaluated using the candidate ranking verification algorithm

Therefore, the results for the latter is simply the degenerate case for the rejection network where the rejection rate is zero percent (see section 5.4). That is, for the test data there is a $79 \%$ success rate and for the training data there is an $83 \%$ success rate. This is an improvement over [Kornai 1997] which yielded $63.3 \%$ success rate on the same database (CEDAR).

## Chapter 7

## Extension: Parallel Implementation

### 7.1 Motivation and Problem Definition

Handwritten word recognition systems should lend themselves to quick training and perform online operation rapidly.

Quick training is desirable because there are typically large amounts of data [Petrowski 1991]. A system should also provide high throughput when it processes fresh inputs online. This is usually more important because users usually have stringent requirements for online performance at the plant site. A mail sorting system may need to process thousands of pieces of mail in a matter of a few hours.

In this chapter, we consider ways to speed up the runtime of the online handwriting recognition system. A simplified version of the system described in Chapter 3 is built using coarse-grain and fine-grain parallelism. It handles the case where an HMM is used to recognize hand printed words. Instead of having oversegmented bitmaps of words, the inputs are sequences of bitmaps of individual characters. By making this simplification we can use the standard Viterbi algorithm instead of the modified Viterbi algorithm discussed in section 3.2. This implementation serves only to demonstrate how the essential parts of a handwritten word recognition engine can be parallelized. In reality, the full power of a Hidden Markov Model (HMM) is not necessary for this problem.

As in the original model, character bitmaps are mapped to the HMM states ${ }^{1}$ using probability distributions arising from a Neural Network.

[^3]The system is implemented using C and MPI (Message Passing Interface) on a Sun HPC (High Performance Computer) cluster and run in both shared memory and distributed memory modes.

Coarse grain parallelism is used to compute the individual character probability distributions ${ }^{1}$. Two processors alternately compute the Neural Network outputs of odd and even numbered character bitmaps.

Fine grain parallelism is used to compute vector "dot products" in parallel for the forward phase of the Viterbi algorithm. The states of the HMM are distributed among processors in almost ${ }^{2}$ uniform blocks. The search for the maximizing state at the end of the forward phase is implemented using a pointer jumping algorithm.

The basic HMM for handwritten word recognition has twenty seven states. In the general case where we could be dealing with cursively handwritten words, the number of states could represent parts of characters and could grow very fast. This implies that for these cases, parallelism would have a greater impact. Data is gathered for timing purposes but since the implementation does not use the extra states, we do not test for correctness for these cases.

### 7.2 The Model

The Hidden Markov Model for hand printed word recognition has twenty seven states. There is a state corresponding to each letter of the alphabet: state $1={ }^{\prime} a^{\prime}$, state $2={ }^{\prime} b$ ', $\ldots$, and state $\mathrm{N}={ }^{\prime} \mathrm{z}$ ' where $\mathrm{N}=26$. The $27^{\text {th }}$ state is a special start state called the $\pi$ state. All paths through the model must start from this state. The figure in the introduction showing the generic HMM is reprinted in Figure 7.1.

1. Recall that these are outputs from the Neural Network classifier.
2. The last processor gets the remainder after the states are equally split among the processors.

There is a link from the $\pi$ state to every other state. The graph with all the states excluding $\pi$ is fully connected. Each link corresponds to the transition probability of going from the source state to the destination state. As usual, the sum of the probabilities of all the outgoing links from a certain node must be equal to one. In keeping with the Markov property, the transition probability of going to some state only depends on which state you are currently at. These probabilities are derived from a lexicon containing the domain of words that our system will be used on.


Figure 7.1: HMM with a pi state and 26 character states

At each non- $\pi$ state, there is a probability distribution $\mathrm{P}\left(\mathrm{o}_{\mathrm{i}} \mid \mathrm{l}_{\mathrm{i}}\right)$ which maps the observations ${ }^{1}$ to the states. We make the assumption that the handprinted characters are separate. This allows us to have identical indices for the observations and their corresponding states.

Given a sequence of character bitmaps, we would like to identify the single best word that would account for the sequence. We therefore need to maximize the probability of a sequence of character bitmaps O given a state sequence $\mathrm{Q}, \mathrm{P}(\mathrm{QlO})$, which is equivalent to maximizing $\mathrm{P}(\mathrm{Q}, \mathrm{O})=\mathrm{P}(\mathrm{Q} \mid \mathrm{O}) \mathrm{P}(\mathrm{O})$ because the sequence of bitmaps stays fixed.

[^4]The probability of a certain bitmap sequence with feature vectors $o_{1}, o_{2}, \ldots, o_{\mathrm{T}}$ given the state sequence $q_{1}, q_{2}, \ldots, q_{T}$ is

$$
P(Q, O)=a_{\pi, q_{1}} P\left(o_{1} \mid q_{1}\right) a_{q_{1}, q_{2}} P\left(o_{2} \mid q_{2}\right) \ldots a_{q_{T-1}, q_{T}} P\left(o_{T} \mid q_{T}\right)
$$

The $P\left(o_{i} \mid q_{i}\right)$ come from a Multilayer Perceptron (MLP) that takes in feature vectors that have been extracted from the character bitmaps.

### 7.3 Coarse grain parallelism

Coarse grain parallelism is achieved by evaluating the outputs of the MLP for different bitmaps in parallel. In the viterbi algorithm, the probability distribution $\mathrm{P}\left(\mathrm{o}_{\mathrm{t}} \mathrm{l}\right)$ is needed only at each step t .

Therefore, these computations can be computed in parallel and the results can be distributed to the processes representing the HMM states when they are needed.


Figure 7.2: Coarse grain parallelism by alternating MLP evaluations

Two processors are assigned to evaluate the MLP. While one is distributing the results for time $t$, the other is computing the values for time $t+1$.

### 7.4 Fine grain parallelism

### 7.4.1 Forward-phase Viterbi

We use fine grain parallelism in computing the matrix-vector "product" for each step $t$ of the algorithm. I repeat the function evaluated at each state at time $t$ at node $i$ :

$$
d_{t+1}(j)=\max _{1 \leq i \leq N} \quad\left\{d_{t} a_{i j} P\left(o_{t+1} \mid j\right)\right\}
$$

This can be done independently for each state i. Therefore, we can introduce parallelism by distributing the $\mathrm{n}=26$ states over p processors.


Figure 7.3: $\mathrm{n}=6$ nodes or states, $\mathrm{p}=3$ processors or processes

1. This operation actually replaces the addition operation in a Euclidean matrix-vector product with a maximizing operation. The analogy is used here because the runtime of both operations are the same.

The Viterbi algorithm now takes up the following amount of time:

$$
c_{1}(p, n)+c_{2}\left(\frac{n}{p}\right)
$$

$c_{1}$ is the cost of broadcasting $O(n / p)$ values to $P$ processors; $c_{2}$ is the cost of doing $n / p$ "dot products"1. The problem is to empirically determine the value of $p$ that minimize $c_{1}(p, n)+c_{2}(n / p)$.

This optimization problem is naively stated because it does not consider the time for backtracking phase of the Viterbi algorithm. If the states are spread across many processors, more communication time is used up querying the right processors for the pointers to the next state in the sequence. Hence, the optimization problem is to minimize the total time taken up by the forward phase and the backtracking phase.

Note that because the amount of time taken by the backward and forward phases are both proportional in the same way to the number of character bitmaps, the minimization problem is independent of word lengths.

### 7.4.2 Pointer jumping algorithm for finding maximum

At the termination of the forward phase of the Viterbi algorithm, we need to find the state that has the maximum probability. To do this, a pointer jumping algorithm was used to convert the $\mathrm{O}(\mathrm{n})$ problem into an $\mathrm{O}(\log \mathrm{n})$ problem [Cormen, Lieserson, Rivest 1990]. However, this code proved to be slow and was later taken out.

[^5]
### 7.4.3 MLP Evaluations

We could also use fine-grain parallelism for evaluating MLPs. The general approach is similar to the one for the Viterbi algorithm: we realize that we can split up matrix-vector multiplications into independent dot products evaluated at each node. This allows us to distribute the nodes among P processors. This approach is not implemented in the current system.

### 7.5 Results

The parallel implementation shows a factor of 2 speedup for HMMs with at least 1500 states peaking at 2 to 4 processors. The coarse grain parallelism for the MLPs generally improved performance and the pointer jumping algorithm did not. We only really gain speedups when the number of HMM states $n$ is greater than 500 .


Figure 7.4: Speedup (y-axis) vs. number of HMM state processors (x-axis)

## Chapter 8

## Conclusions and Future Work

### 8.1 Conclusions

A system for offline cursive handwritten word recognition was built. As the framework for the recognition task, a variant of VSLHMM was introduced: HMM with Multiple Observation Sequences (HMMMOS). The results show a close visual correlation between the hypotheses and the actual word bitmaps. When global context and a lexicon are used to verify and rank hypotheses, the system recognizes words better than other systems when applied to the Center for Excellence in Document Analysis and Recognition (CEDAR) database.

Among the three neural classifiers tried, only the Multilayer Perceptron (MLP) performed satisfactorily. This might be explained by the fact that MLPs have been successfully used in the past for this kind of application, while Hierarchical Mixture of Experts (HME) and Recurrent Networks have been relatively unexplored.

Another MLP was successfully trained and used to reject bad inputs in order to increase the recognition rate.

Finally, a parallel implementation for a simpler system shows that adding one to three processors doubles the online recognition speed.

### 8.2 Future Work

8.2.1 Compare the three neural classifiers.

It is interesting to make the HME with Generalized Linear Model (GLIM) experts converge or find out why it does not converge. In [Waterhouse 1993] an HME with GLIM
(Multinomial) experts is applied to voice recognition. Although this did not converge either, a setup that did converge in that thesis is one that had a tree of HMEs, each with only two logistic functions. This kind of architecture is usually undesirable because the output probabilities are not normalized and there is no reliable way to compare them to even get a ranking. However, it would be interesting to see if this alternative architecture converges for our system.

Because the neural classifier is the basis and bottleneck for the HMM, it is important that we get good performance from it. It is therefore of interest to compare the three neural classifiers tried here and possibly others.

### 8.2.2 Stochastic Context Free Grammars

It would be interesting to implement and compare an SCFG system with the HMMMOS. As mentioned in section 1.6, the two are similar and a further study to try to reconcile them might lead to better performance.

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[^0]:    1. This is the state from which the arrow representing the directed edge starts. The destination state is where the arrow points.
[^1]:    1. To keep the algorithm running, at least one candidate is put in the list even if all the candidates do not exceed the threshold probability.
[^2]:    1. These features are discussed in section 5.3.
[^3]:    1. The HMM states correspond to the letters of the alphabet ' $a$ ' through ' $z$ '.
[^4]:    1. These are the feature vectors that are derived from the character bitmaps.
[^5]:    1. Again, it is not really a dot product because we compute the maximum instead of the sum.
