

THE SYNTHESIS OF SPEECH
USING A DIGITAL COMPUTER

by

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ABSTRACT

A program has been written for the IBM-704 digital computer to synthesize speech. The input requires detailed specification of the speech spectrum as a function of time. The output is audible samples of 2.4 seconds duration with a signal-to-noise ratio up to 36 db and a frequency response up to 7 kc. Included is a complete SAP listing and flow-charts of the program. The synthesis of each sample takes from 6 to 10 seconds. The IBM-704 "Direct-Data" attachment and a digital-to-analog converter is required.

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Introduction

The purpose of this work is the development of a device to synthesize speech^{which} will be able to manipulate the acoustic "building blocks" of speech in a very flexible manner. Today a synthesizer is primarily useful for studying speech by testing various hypotheses about the interactions of various speech sounds and the nature of our perception of speech. Eventually it is hoped that speech may be added to our methods of communicating with machines. Before this can be done we need to have a fairly detailed knowledge of the acoustics of speech.

Acoustics of the Vocal Tract

The vocal tract begins at the vocal cords in the larynx and then divides, one section terminating at the mouth approximately 17 cm. away and the other section terminating at the nose a slightly greater distance away as shown in Fig. 1.

For a large class of speech sounds we need only consider the oral cavity, as the nasal cavity is virtually isolated from the system. Thus in essence the vocal tract can be considered to

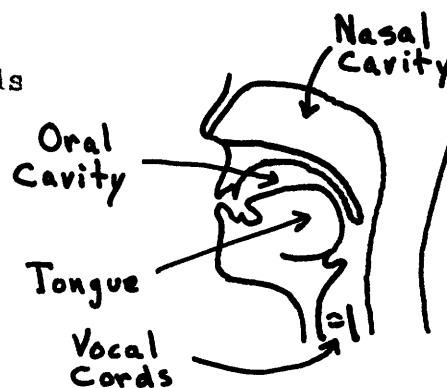


Fig.1

be a tube of fixed length and variable cross-sectional area. The area varies from 0 to 15 cm.².

Electrical Analogy

If we define an acoustical impedance as the ratio of the applied pressure to the volume velocity ($\text{cm.}^3/\text{sec.}$), we find that the vocal tract is analogous to an electrical transmission line with distributed inductance $L=p/S$ and capacitance $C=S/c^2p$ per unit length. Here p is the density of air, c is the velocity of sound and S is the cross-sectional area of the tract at the corresponding point.

A transmission line may be simulated by lumped elements, inductors and capacitors, provided the wavelength λ is long compared to the section being represented. For speech the highest frequency of normal interest is of the order of 6 kc. with a corresponding wavelength of $\lambda=c/f$ which is $3 \times 10^2 / 6 \times 10^4 = 0.05$ m. or 5 cm. Thus the vocal tract can be adequately represented by approximately 15 sections. In this analogue the pressure in the system corresponds to the voltage and the current to the volume velocity. The source (for vowels) is represented by a high impedance (current) generator. The output is taken as the pressure (voltage) at the lips which are loaded by a dissipative impedance representing the loss of energy to the surrounding air. See Fig. 2 below.

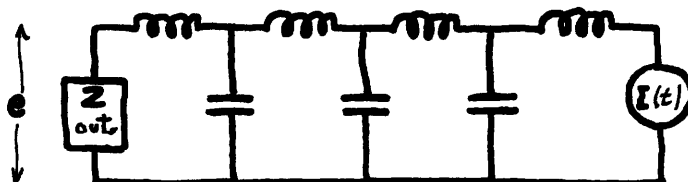


Fig. 2

Solution of the Circuit Equations

The elements in the circuit obey the simple equations:

$$i = C \frac{de}{dt}, \quad e = L \frac{di}{dt}, \quad e = iR.$$

By ordinary circuit analysis we obtain a differential equation of the form:

$$b_n \frac{d^n e}{dt^n} + b_{n-1} \frac{d^{n-1} e}{dt^{n-1}} + \dots + b_0 e = i$$

We are primarily interested in the homogeneous or force-free ($i=0$) solution. The standard method for solving the above linear, homogeneous, differential equation with constant coefficients is to make the substitution $e = e^{st}$.

Doing this we obtain the algebraic equation:

$$b_n s^n + b_{n-1} s^{n-1} + \dots + b_0 = 0$$

Denoting the roots of this equation by s_n we have the homogeneous solution:

$$e = B_n e^{s_n t} + B_{n-1} e^{s_{n-1} t} + \dots + B_0$$

where the values of the constants, B_n , are determined by the initial conditions.

The actual system has distributed losses due to viscosity, cavity wall absorption and due to the fact that the process is not adiabatic as the calculations suppose. In the case of losses the roots are in general complex numbers, $s_i = \sigma_i + j\omega_i$, and the solution for the force-free condition can be written in the form:

$$e = B_n e^{\sigma_n t} \sin(\omega_n t + \alpha_n) + B_{n-1} e^{\sigma_{n-1} t} \sin(\omega_{n-1} t + \alpha_{n-1}) + \dots$$

The Production of Vowels and Consonants

Vowels: The current source can be thought of as delivering impulses spaced at intervals of $1/f_0$ where f_0 is the fundamental pitch or voicing frequency. The impulse determines the initial conditions and for the remainder of the interval the circuit decays by the force-free solution above. See Fig. 3 below.

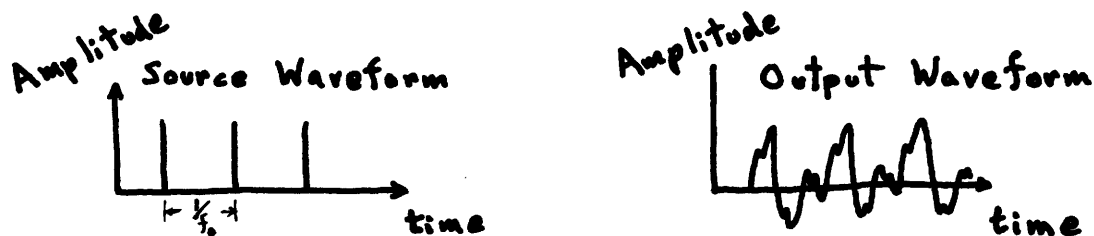


Fig. 3

In general the acoustic network acts as a filter on whatever source spectrum is present, as shown in Fig. 4 below.

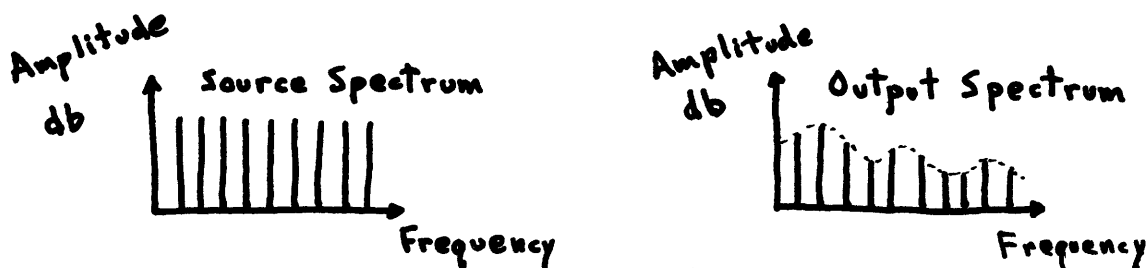


Fig. 4

Consonants: The production of a fricative consonant such as (s) proceeds in this manner: A constriction, in this case caused by the tongue pressing against the front of the roof of the mouth, causes a high velocity, turbulent flow of air. This is a source of noise which is filtered by the cavities ahead of it, around the teeth and to some extent by those back of it, giving a spectrum of the type shown in Fig. 5 on the next page.

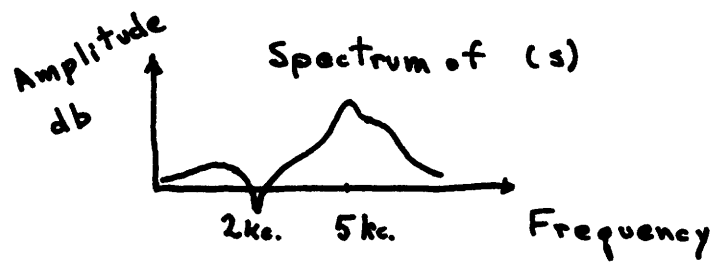


Fig. 5

The equivalent circuit for the production of fricative consonants seems to involve a noise voltage source at the place of the constriction as shown in Fig. 6 below.

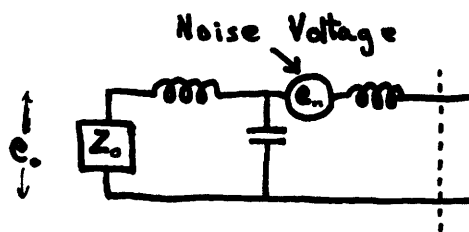


Fig. 6

Previous Methods of Synthesis

For the synthesis of speech a number of methods have been used. Probably the two most flexible ones are those at MIT^{16*} and the Haskins Laboratory.⁴ The first uses a dynamic analog vocal tract. It is an electrical analogy in which the "cross-section" of the tract is continuously adjustable by purely electrical means (saturable reactors and reactance tube circuits.) The other at the Haskins Laboratory in New York is a device allowing playback of hand drawn speech patterns. The patterns are drawn to any desired complexity and the sound is reproduced by the playback device.

A number of other devices have been built^{9,10,11,13,14,15,} but they all suffer in being difficult to "program" for complex or rapidly changing speech sounds.

*Superscripts refer to the bibliography on pgs. 16 - 18 .

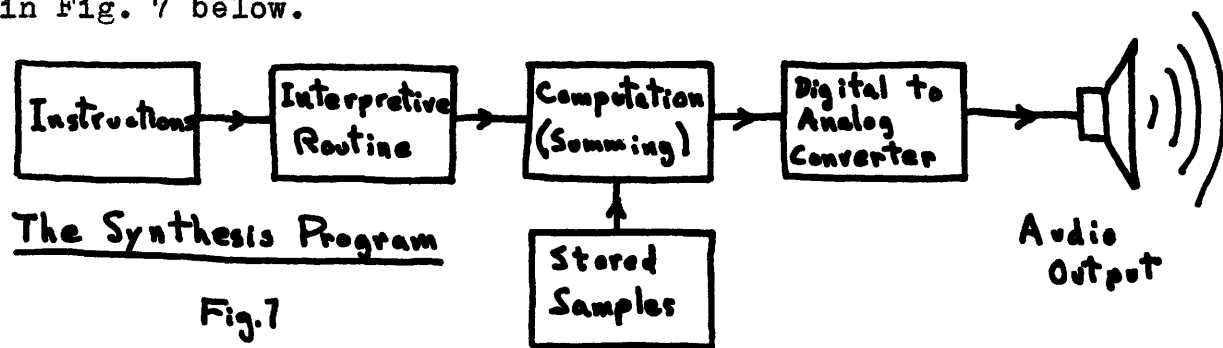
DIGITAL COMPUTER METHODS

Introduction

It is possible with certain high speed digital computers both to compute the waveforms and to have the waveforms presented at a loudspeaker in their normal audible form, i.e., as speech sounds. The computer can also manipulate at high speed the instructions specifying the nature of the waveforms.

The differential equation describing the vocal tract could be converted to difference form and solved continuously to yield the desired output. However, this involves rather high order differences and is thus time-consuming. It is intriguing because it can be made quite exact and still retains the analogue form fairly well. A faster approach, and the one utilized, is to compute a number of samples of solutions and have the program simply extract the desired sections of the solutions from its table and present them at the output. The stored solutions would contain damped sine waves of many different frequencies and filtered noise of various bandwidths and frequencies. This produces a quazi-static solution quite similar to the WKB solution of a wave equation. This approach is justified by the fact that the vocal tract parameters change much more slowly than the characteristic decay rates of the cavity responses.

A program has been written by myself to perform these operations on the MIT-IBM 704 computer. The program has been designed primarily for high speed and flexible operation. An outline of the program is shown in Fig. 7 below.



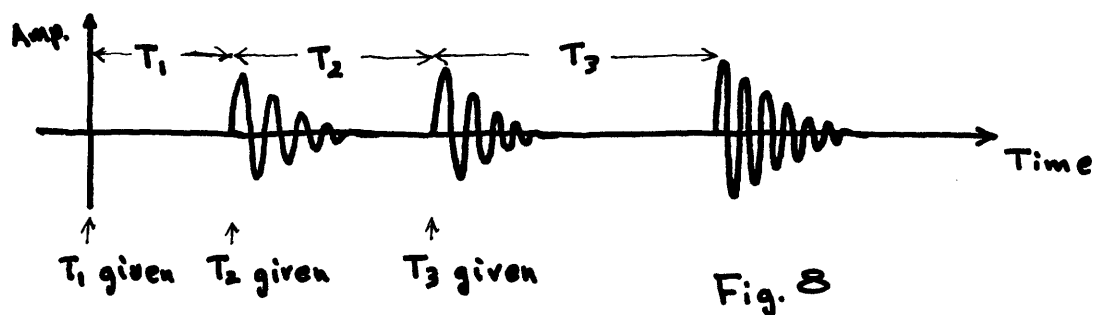
Digital to Analog Conversion

The conversion of the synthesized samples in their digital form to an analog voltage which can be recorded or reproduced in audible form is simply accomplished by a converter such as the one shown in reference 26, pg. 530. Within the machine samples are computed to six binary places of accuracy. This gives a maximum of 36 db signal-to-quantization level ratio, a rough measure of the maximum signal-to-noise ratio. Frequency response is limited by the maximum obtainable read-out rate of 14,000 points per second. The highest component that can be transmitted is 7,000 cps.; all higher are reflected into the 0 to 7,000 cps. region and appear as noise.

The Instructions

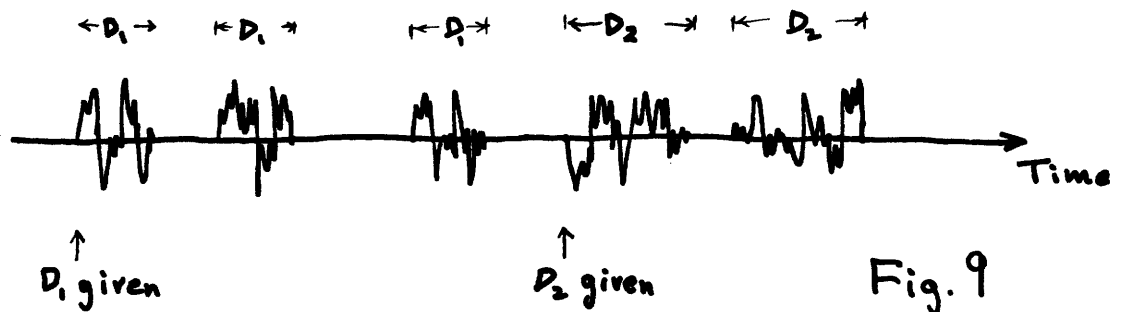
When summing a number of sine and noise samples, the instructions must specify the nature of the samples (frequency, amplitude, phase, bandwidth, etc.), the point at which the sample is to begin, and finally the duration of the sample. The four instructions used to specify this information are: "Time", "Duration", "Formant," and "Noise."

Time A marker is kept within the machine which denotes the present point of time. This point is a certain location in the 704 magnetic core memory. When any sample is specified it will begin at this point. When the Time instruction is given (abbreviated by T) the marker is moved ahead by the amount specified by the instruction. Thus a series of T instructions interspersed with other instructions might give the output shown in Fig. 8 below.



Duration It is sometimes desired to limit the duration of a sample. Often a noise burst is desired which is much shorter than the particular sample in storage. By giving a duration instruction (abbreviated by D) we can accomplish this limiting. The amplitude of a noise sample is set and

cannot be changed while it is being summed. Therefore, if a burst of increasing amplitude is desired, a series of short bursts of length limited by D can be specified. Each of these bursts would have a greater amplitude than the preceding one. The value of the duration is always that of the last given; it need not be stated frequently. The effect of the D instruction is shown in Fig. 9 below.



Formant The Formant instruction specifies a single damped sine wave. This instruction (abbreviated F) specifies the phase (as plus or minus), the frequency, and the amplitude. There are normally about 50 frequencies available between 100 and 5,000 cps. The amplitude is adjustable in 3 db steps over a range of 36 db.

Noise The Noise instruction (abbreviated N) specifies the amplitude in 6 db steps and the frequency of the noise. The exact character of the noise samples is yet to be decided.

Summation

It is normally necessary to sum a number of different damped sines or noise. This is done very simply by letting them overlap or coincide in time, when specifying them in the instructions. Thus, if a number of different F instructions are given without an intervening T instruction the sum of the various F's will be formed, similarly for noise. There are no restrictions on the ordering of the instructions -- they are completely independent in that sense. To produce a section of a vowel we might give a sequence of the form FFFTFFFTFFF..... which would produce a sum as shown in Fig. 10 below.

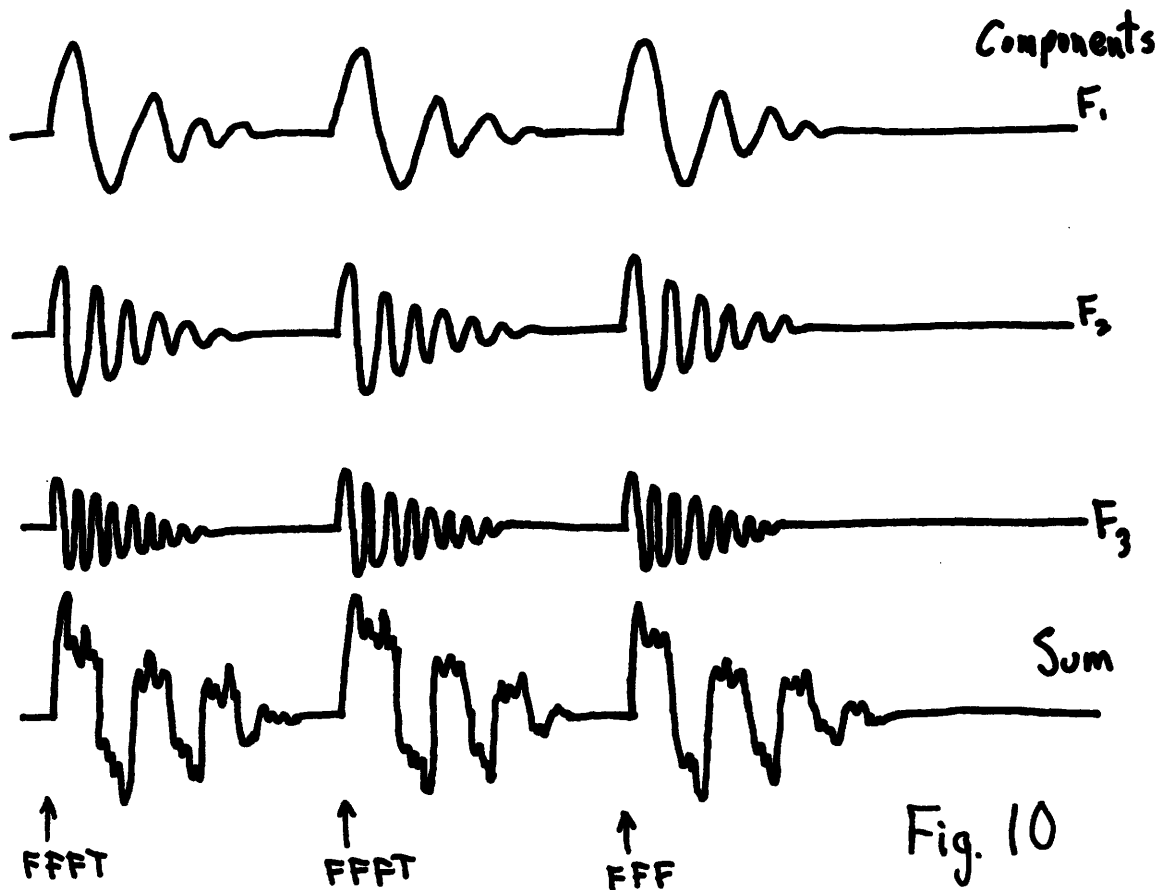


Fig. 10

Preparation of Instructions

The instructions are punched on IBM cards to be read into the machine. Once a large number of cards is punched they need not be re-punched, they need only be re-ordered to form new sounds. If a number of identical cards are required they can simply be reproduced from a "master" (cf. the last 15 instructions in Appendix F.)

The program considers the instructions in order, one-by-one. It is designed so that the parameters such as the frequencies and the bandwidths of the samples can easily be altered. The noise samples are actually computed by the synthesis program itself using random numbers, and the T, D, and F instructions. The noise can thus be made to have almost any spectral characteristics desired. The program is also designed to be readily modified to receive new instructions that might be wanted to store the samples on magnetic tape, display the samples on an oscillograph screen or present them at the audio output for recording.

The Piece-wise-linear Program

The primary disadvantage of the program outlined above is that we must write too many instructions for each sample we wish synthesized. Often there is a great deal of redundancy in these many detailed instructions because the sounds are changing slowly or in some very regular way, e.g., the fundamental pitch or some of the

formants may be changing approximately linearly with time during some interval. For cases such as these we use a coding scheme to eliminate a great deal of this redundancy. Even simple schemes can reduce the number of instructions necessary for a 2 second sample from many hundred to easily less than one hundred. The method of doing this is to write a second program which accepts the more concise instructions and translates them into a series of many more of the T, D, F, and N instructions. This operation is outlined in Fig. 11 below.

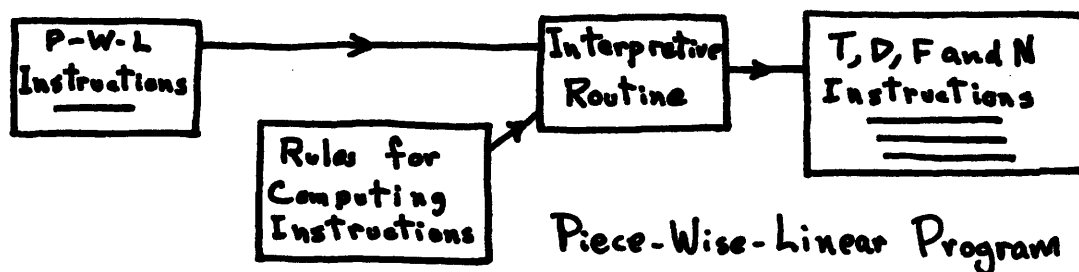


Fig. 11

At this point it is useful to discuss a convenient representation of the speech we want to synthesize.

A common way to represent the dynamic acoustical properties of speech is by "visible speech" patterns²¹. These are produced by certain devices used for the analysis of speech. The patterns are constructed so that the vertical axis represents frequency, the horizontal time, and the density of the shading represents intensity. Vowels then appear as horizontal bars called formants and fricatives appear as grey regions of no very definite structure in the high frequency region. This is shown in Fig. 12 at the top of the next page.

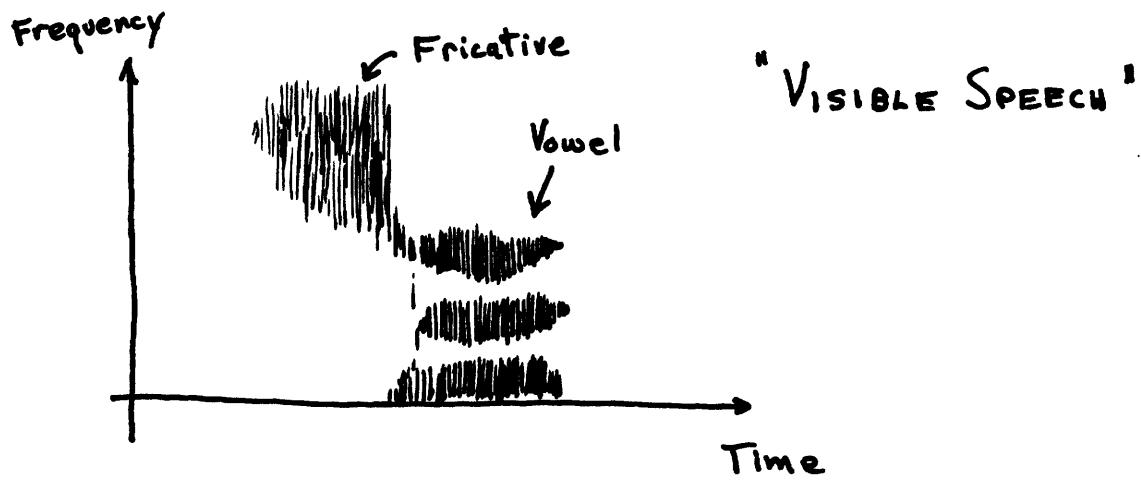


Fig. 12

To code for the piece-wise-linear program we first approximate the movements of the formants by straight-line segments. At each break-point we state the properties of the formant somewhat as an F instruction does. We do the same for the noise and for the fundamental frequency. This is shown in Fig. 13 below.

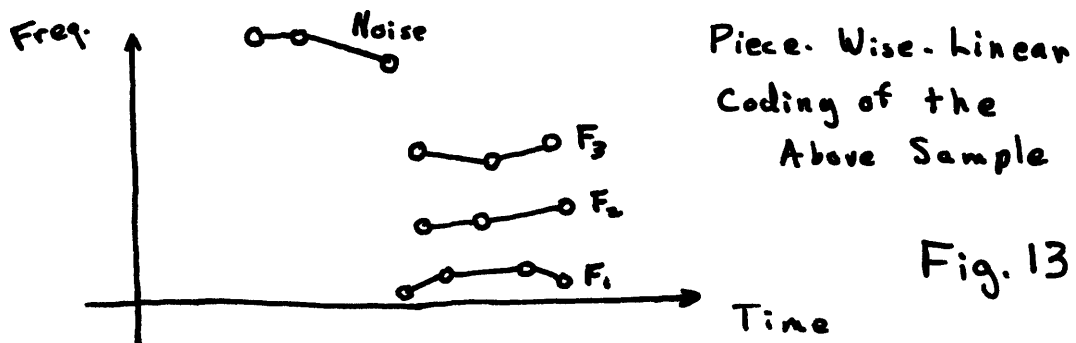


Fig. 13

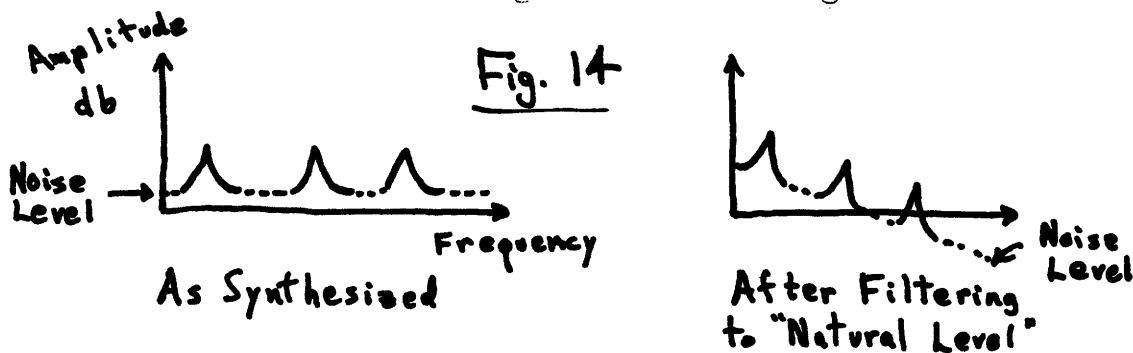
These patterns are then coded and presented to the piece-wise-linear program which interprets them and composes the proper detailed instructions.

The present writer has not worked on the above program in detail but another student is studying the problem.

Pre-emphasis

By programming unnaturally large amplitudes for the higher frequencies we can, in an approximate way, pre-emphasize them. Filtering the output will bring

them back to their "natural" level but now with less noise in the higher frequencies. This is done, of course, at the expense of adding noise and reducing the dynamic range in the lower frequency region. It has been found that the higher frequencies are quite important in the production and recognition of speech even though their normal amplitudes are lower than the low frequency region. The action of the noise and filtering is shown in Fig. 14 below.



Program Speed, Testing, and Future Plans

The maximum length for a single continuous sample that can be read out of the machine is 2.4 seconds at 14 kc.

In synthesizing a sample involving the first three formants the program can run as fast as twice real time, i.e., it would take the machine 1.2 seconds to synthesize a sample which would last 2.4 seconds. The slowest rate at which synthesis proceeds is half of real time. The major delay in time is preparing the sample for the output device which takes half of real time. Thus it would normally take between 6 and 10 seconds to compute the sample and ready it for read-out.

The program has been tested and debugged to a large extent. A cursory examination of the scope display of the synthesis using the instructions in Appendix F indicated that the program was functioning properly though no films have been received as yet. Further testing using scope display is in order.

The next steps in the development of this program might be the inclusion of tape routines for storing the bank samples, the piece-wise-linear program, and finally the setting up and use of the digital-to-analog equipment. One of the runs of the program should be devoted to computing the noise samples.

A reasonable arrangement for the actual operation of the program would be to have all of the programs and bank samples on tape. Then the piece-wise-linear instructions would be read in from the on-line card reader. Synthesis would then proceed, the samples finally being recorded on an audio tape recorder for listening tests.

* * *

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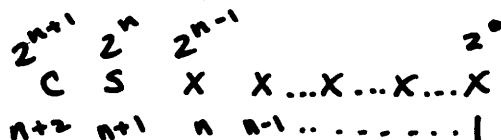
Appendix A

Sub-word Arithmetic

When little accuracy is desired it is useful to divide a 36 binary digit 704 word into a number of smaller words called sub-words. Each sub-word carries a sign with it and one position must be left open for a carry so that one sub-word will not interfere with another. The sign and carry bit are indicated by S and C respectively. The negative is handled by using the two's complement. These numbers are convenient because addition, subtraction and multiplication can be performed on all of the subwords in a 704 word without ever "unpacking" them. In this program five sub-words are stored in a 704 word as follows:

1	2	3	4-8	9	10	11-15	16	17	18-22	23	24	25-29	30	31	32-36
x	C	S	...	C	S	C	S	C	S	C	S
			5			4			3			2			1

A single n-place signed sub-word is of the following form:



We have assumed the binary point to be to the right of the first place but the argument below can simply be generalized to this case (normally the sub-word is surrounded by others.)

The n+1st place designate the sign, S. A zero in S indicates a positive number, a one indicates a negative number.

The $n+2$ nd place is a carry bit for absorbing any changes of sign -- it is erased after each computation. If a number of these sub-words in a 704 word are added to another similar 704 word, each of the n -place numbers will be added algebraically to the corresponding number (sub-word) in the other 704 word.

The negative of a number is obtained by subtracting it from 2^{n+1} (complementing it.) (and then adding 1)

First we restrict all numbers of interest so that

$$0 \leq a, b < 2^n \quad \text{and} \quad |a| + |b| < 2^n$$

We will show that the negative of a number always has a one in S. Trivially a number $< 2^n$ has a zero in S so that complementing it always produces a one in S.

It is interesting to note that the number 2^n has the properties $(-2^n) = 2^n$ and $2^n + 2^n = 0$.

Laws of Addition

$(+a) + (+b) = (a+b)$ by the normal addition process.

For $(+a) + (-b) = a-b$ we have $a > b$ or $2^n > a-b > 0$ so that $2^{n+1} + 2^n > (+a) + (-b) > 2^{n+1}$ giving a zero in S, a positive number as we had assumed.

If $a = b$, $(+a) + (-b) = 2^{n+1}$, a zero.

If $a < b$ then $0 > a-b > -2^n$ or $2^{n+1} > (+a) + (-b) > 2^{n+1} - 2^n = 2^n$,

giving a one in S.

For $(-a) + (-b)$ we have $2^{n+2} > (-a) + (-b) > 2^{n+2} - 2^n$,

giving a minus sign in S.

All of the sums considered here are $< 2^{n+2}$ meaning that no C ever overflows into and interferes with the sub-word on the left.

Subtraction of the sub-words "in bulk" is similar to addition but the C bits of the minuend must be loaded with 1's to prevent borrowing from the sub-word on the left.

Multiplication

We can multiply a number by a factor of 2^{-m} , if it is positive, by shifting it m places to the right and inserting zeros on the left. If we consider the digits shifted past position 1 to be lost we have, originally,

$$c = k_{n-1} 2^{n-1} + k_{n-2} 2^{n-2} + \dots + k_0 2^0$$

After shifting m places we have decreased each exponent by m , giving,

$$c' = 2^{-m} c \quad \text{to within } 2^{-n} = 1.$$

If c is the negative of a number, $c = 2^{n+1} - a$, we insert 1's on the left after shifting.

Originally $\underset{c}{0} \underset{s}{1} x \dots x$ is shifted to $\underset{c}{00} \dots \underset{s}{0} x \dots x$,

and we add to this $\underset{c}{0} \overset{1}{\dots} \underset{s}{1} 0 \dots 0$.

The number added is $\underset{c}{2^m} (2^{n+1} - 1) - (2^{n-m+1} - 1) = 2^{n+1} - 2^{n-m+1}$

so that $c' = 2^{-m} (2^{n+1} - a) + 2^{n+1} - 2^{n-m+1} = 2^{n+1} - 2^{-m} a$

which is the negative of the number $2^{-m} a$ as desired.

If a digital-to-analog converter is used which only handles absolute values we add 2^n to each of the sub-words. This is a 704 word with 1's in the S bits and 0's elsewhere. We can also accomplish the same operation by simply changing the sign which can be seen to be equivalent.

Appendix B The 704 Program in Detail

The 704 program as tested in May 1959 is made of two basic sections. First is the synthesis program itself, and second the auxiliary routines -- one to initially compute the damped sines and another to display synthesized waveforms on the cathode ray tube unit attached to the computer.

Synthesis Program

The present point of time may occur at a sub-word within a 704 word so that addition, when it occurs, will involve shifting the stored samples to the right or left so that their beginning will correspond with the present point of time. On the other hand the present point of time may be at the beginning of a 704 word which means no shifting is required. Since the second operation is much simpler than the first, an optimum program could be written that would make the simple routine run much faster. Also, multiplication may or may not be required and there is quite a difference of speed when multiplication is omitted.. To take advantage of these possibilities to gain speed, routines have been designed to handle each case separately. The phase (plus or minus) is also handled separately at times. There are three working areas in the machine: the bank or sample block containing the damped sines and noise, a short work space called the multiply block, and finally the assembly block in which the entire synthesized sample is formed.

The various possible sequences of operations are shown in Fig. 15 below.

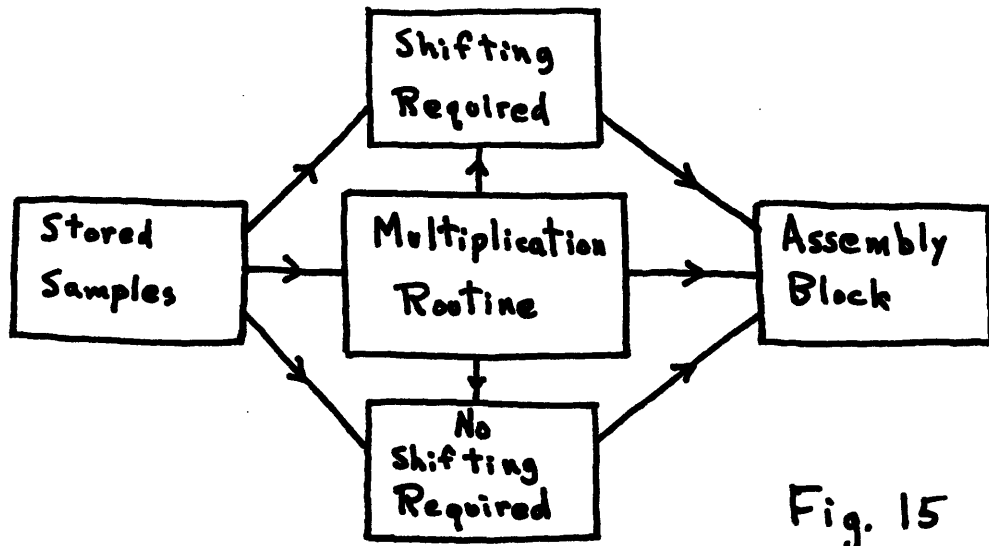


Fig. 15

The instructions are tested by the program one at a time and completely executed before the next instruction is studied. Whenever the synthesis (sometimes called "assembly") program is entered all of the markers are initialized so that it begins at the origin of the assembly block. After an instruction has been executed a return is made to obtain the next instruction, or an error count is stepped by one if an illegal instruction was given (amplitude or frequency out of range), or the assembly may terminate. Termination occurs if an all zero instruction is found (this is a T zero instruction,) or if the present point of time becomes located near the end of the assembly block.

The T instruction causes the setting of three markers: the present point of time measured in sub-words from the beginning of the assembly block, the number of 704 registers from the beginning of the assembly block, and the position of the present point of time in sub-words as measured from the beginning of the 704 word which contains it.

The D instruction sets the duration as measured in sub-words. If the duration is found to be greater than the length of the longest sample (a noise sample), then it is reduced to this value.

The F instruction first computes the phase of the sample as plus or minus and sets a marker denoting this. Since the multiplication routine can only handle shifts of multiples of 6 db in amplitude (powers of 2), and 3 db steps are desired, each sine is stored at two amplitudes separated by three db. Thus both the amplitude and the frequency enter into the calculation of the location in storage. Both the number of places to shift, for multiplication, and the sample location are then computed. Finally, advantage is taken of the fact that the sine waves damp out in shorter times for lower amplitudes. If this damping time is shorter than the duration time already specified, the routines handle the sample for the minimum time necessary and do not compute useless strings of zero's. The F routine then transfers control to the proper assembly routine for the actual synthesis.

The N instruction is quite similar to the F except that the phase is always plus and because of storage requirements the amplitude is in 6 db steps.

In this program the end of assembly is followed by the (TV) routine for scope display. In a production version the end of the assembly would cause the sample to be stored or to be read out. Then a new set of instructions would be obtained and the next sample synthesized.

Auxiliary Routines

To compute the damped sine waves, one cycle of a 55 cps. tone and one exponential of somewhat longer duration are computed. The nth harmonic is then formed by indexing through the block of the sine wave, picking every nth point. The computations for the indexing are done modulo the block length (which is conveniently $256=2^8$), so that the point picked is always constrained to lie within the one cycle block. This is a very fast way to produce multiples of a waveform in time. These harmonics are multiplied by the decaying exponential, scaled to two different values and the resulting samples stored. All harmonics of 55 cps. up to a pre-set limit are computed and then every even harmonic up to a second limit. The limits in the present version of the program are 550 and 5500 cps, respectively. This routine is finished at a point called "End Computation." At this point a routine can be inserted to store the samples on tape so that they need not be computed for each run.

The scope routine first titles the run and then takes blocks of 100 words (500 sample points) and displays them a number of times, numbering each block sequentially from 1. The right edge, left edge and the x-axis are also displayed. The role of this routine is quite flexible, as it is only needed when there is some question of whether or not synthesis is proceeding correctly -- it is quite useful in debugging.

Appendix C

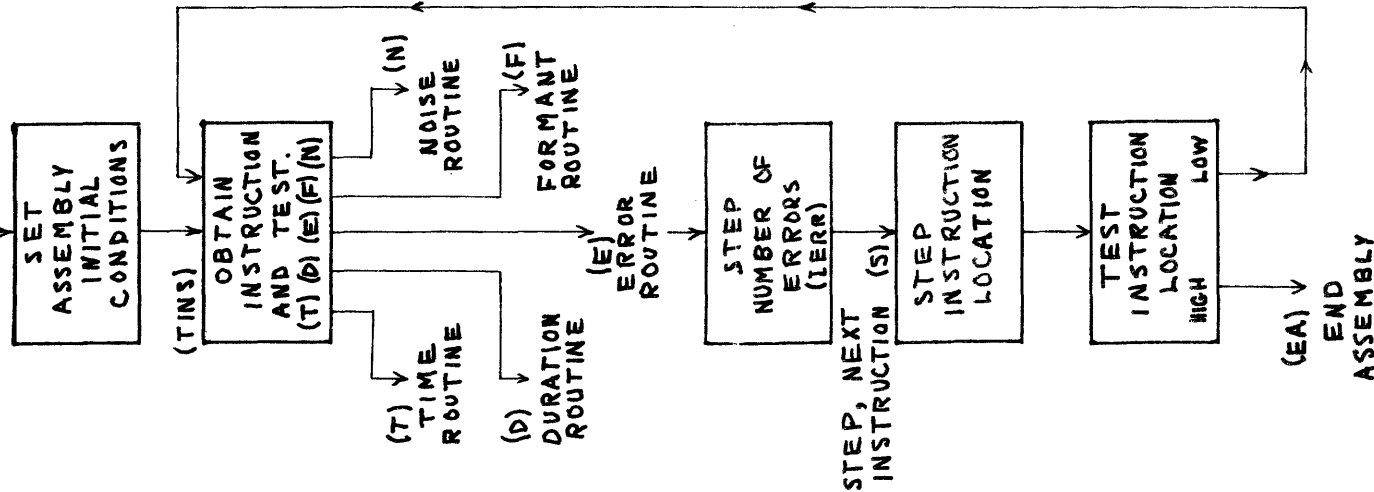
FLOW - CHARTS

General Synthesis Program page 28

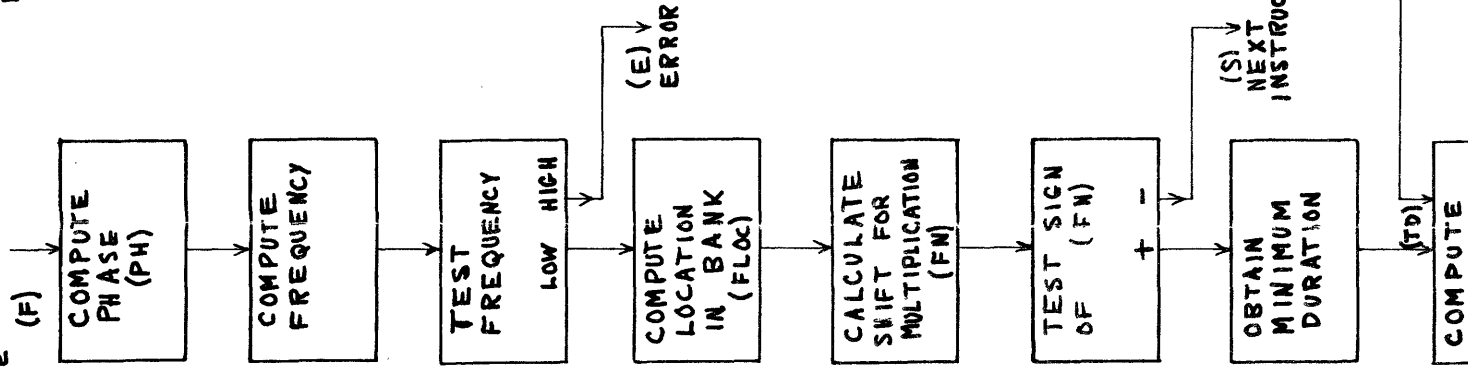
Auxiliary Routines. page 29

On the flow-charts the symbols in parentheses refer to points of entry and data words in the SAP program. The actual locations of the most important of these in the SAP program are indexed on page 30.

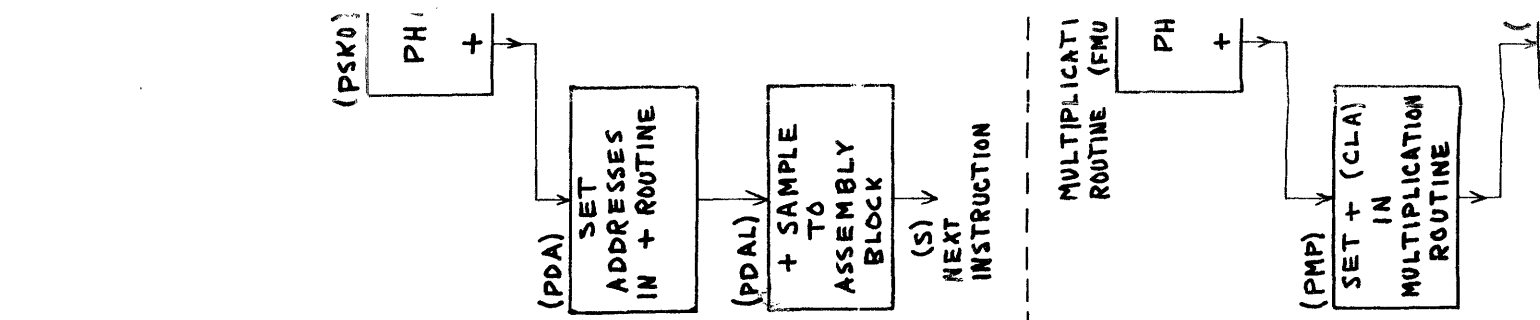
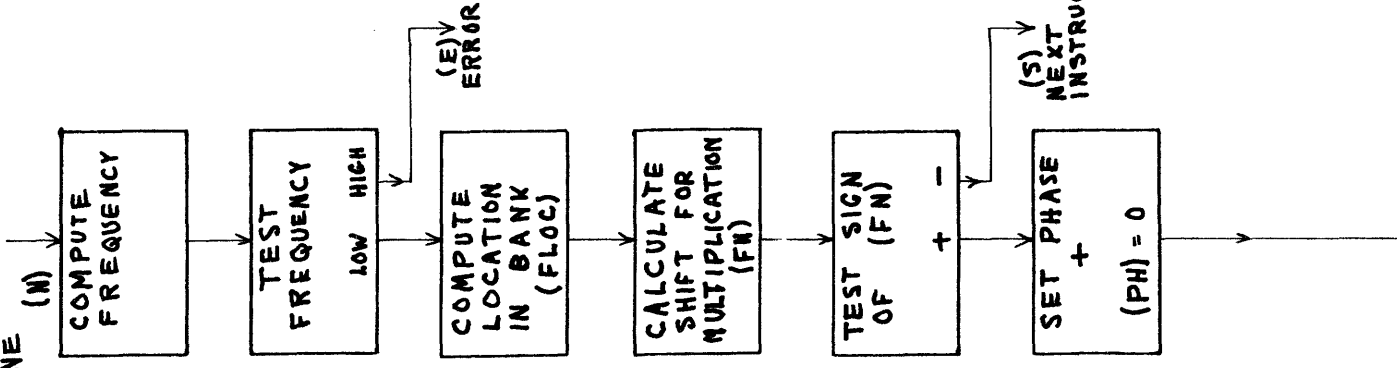
ASSEMBLY ROUTINE (A)



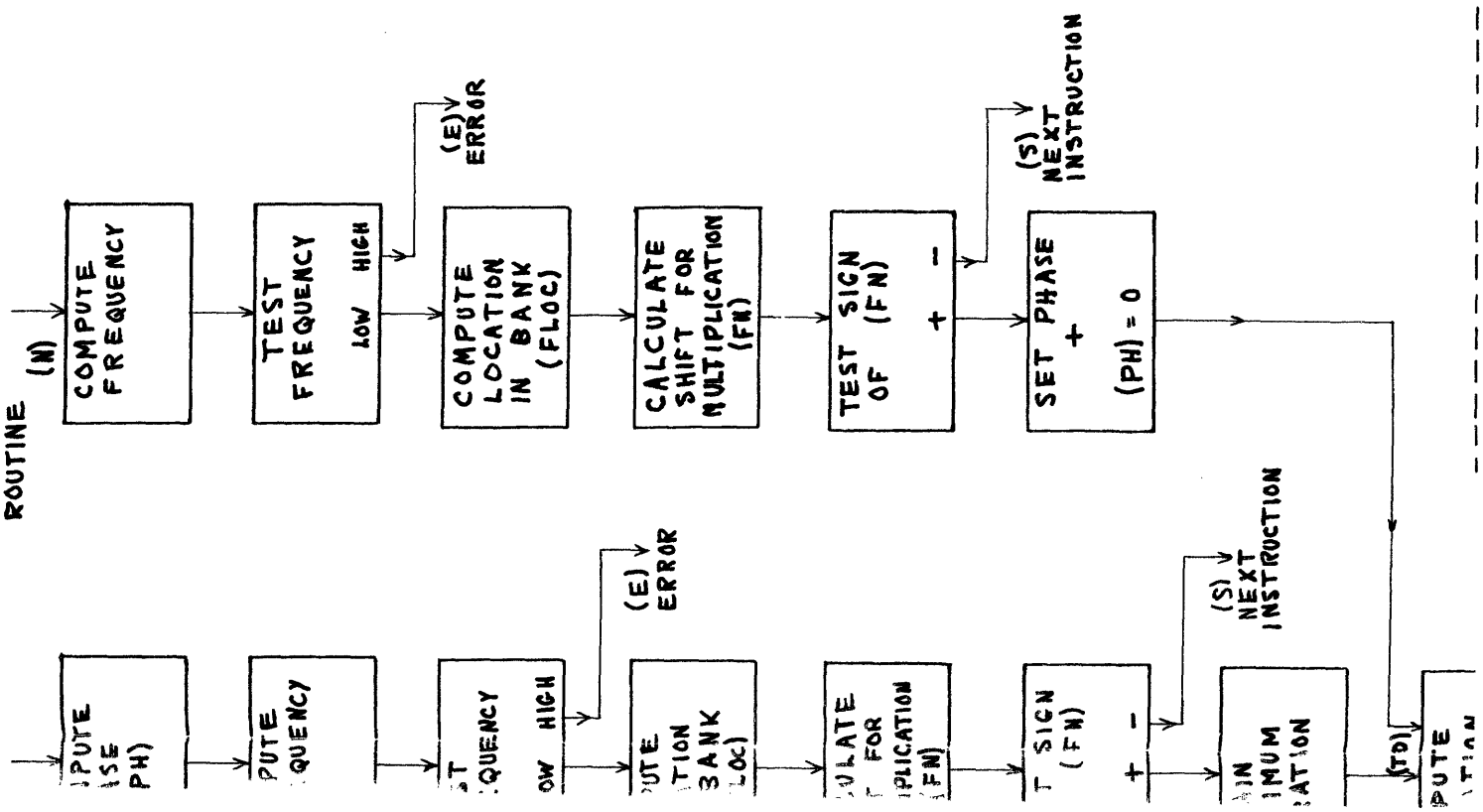
FORMANT ROUTINE (F)



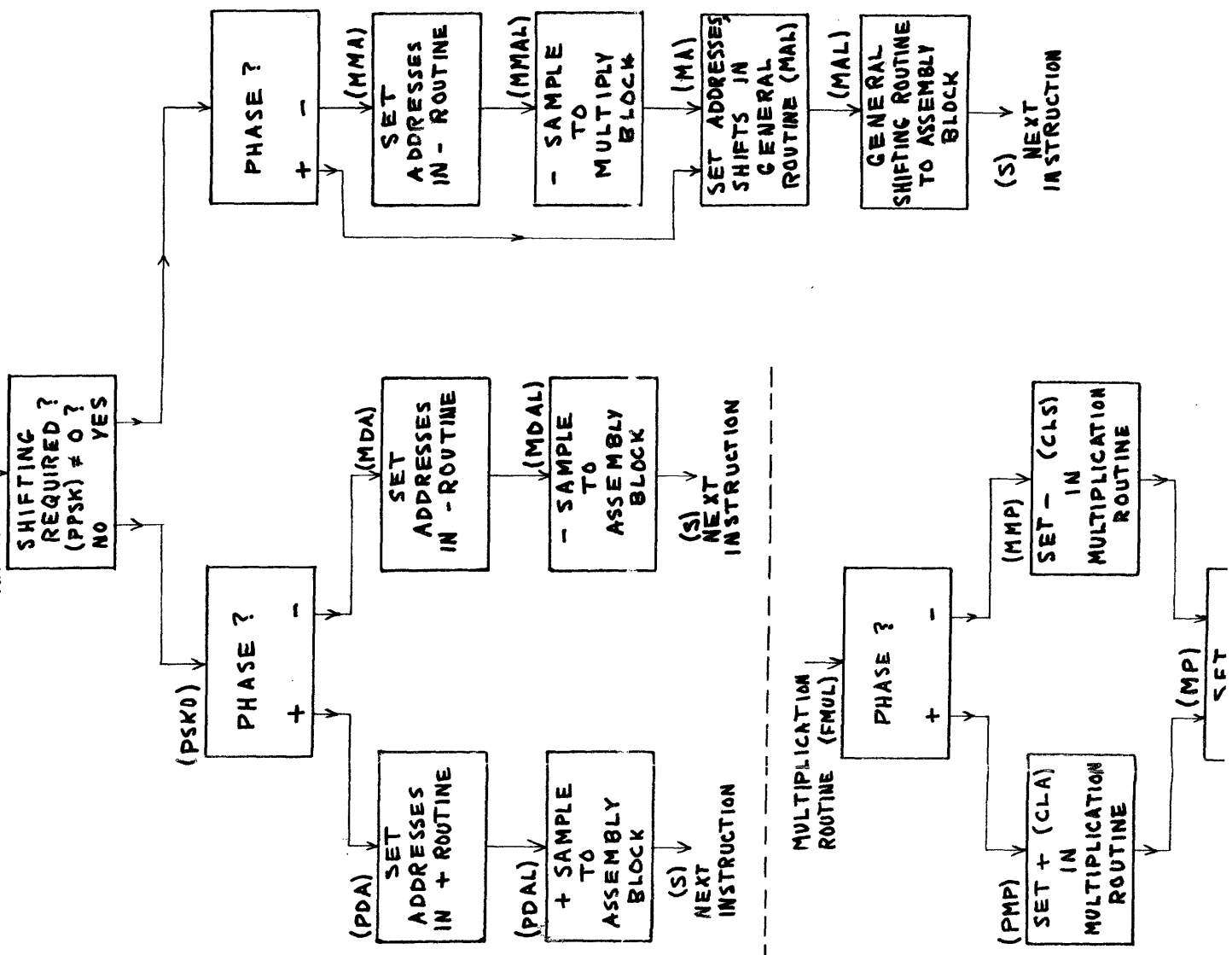
NOISE ROUTINE (N)

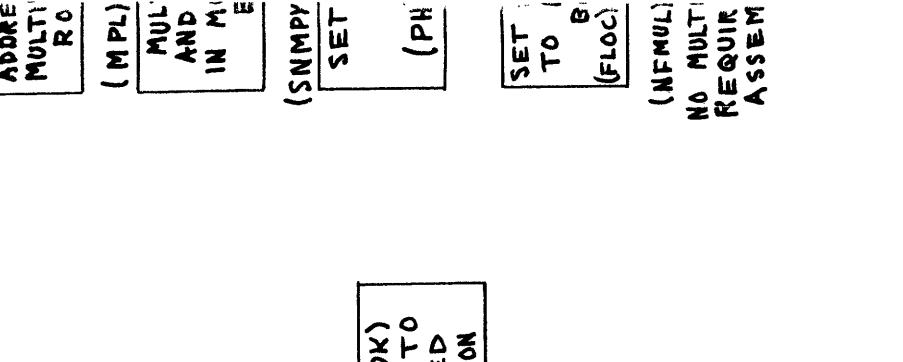
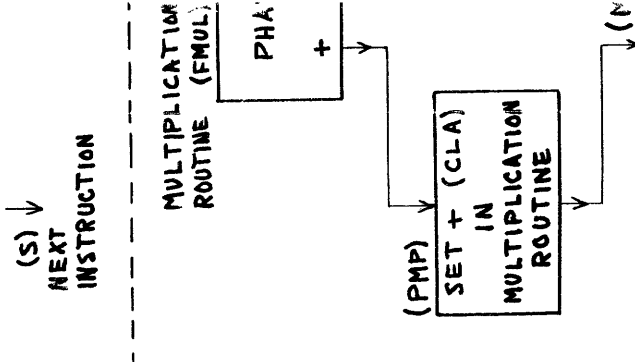
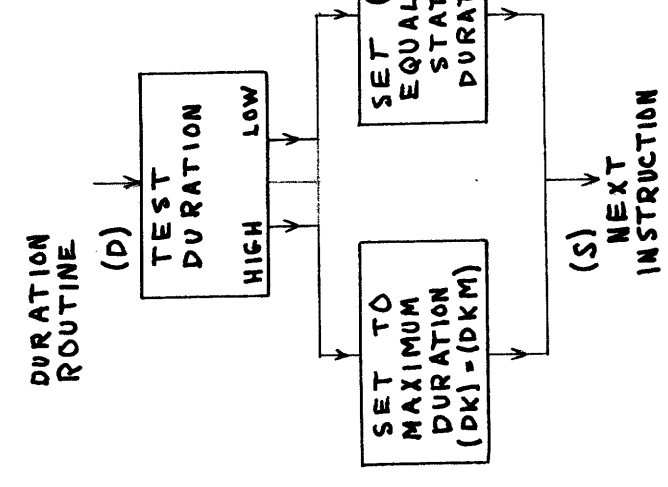
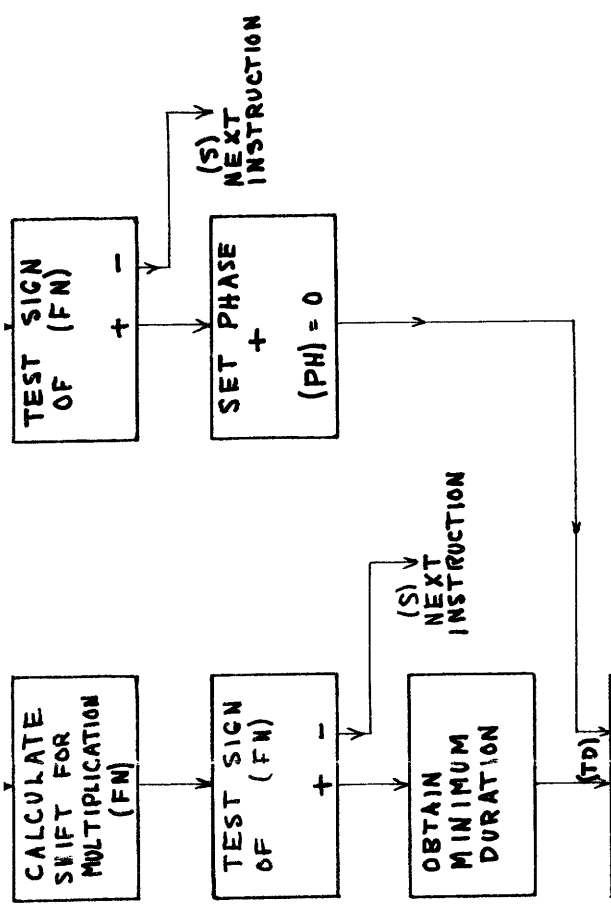
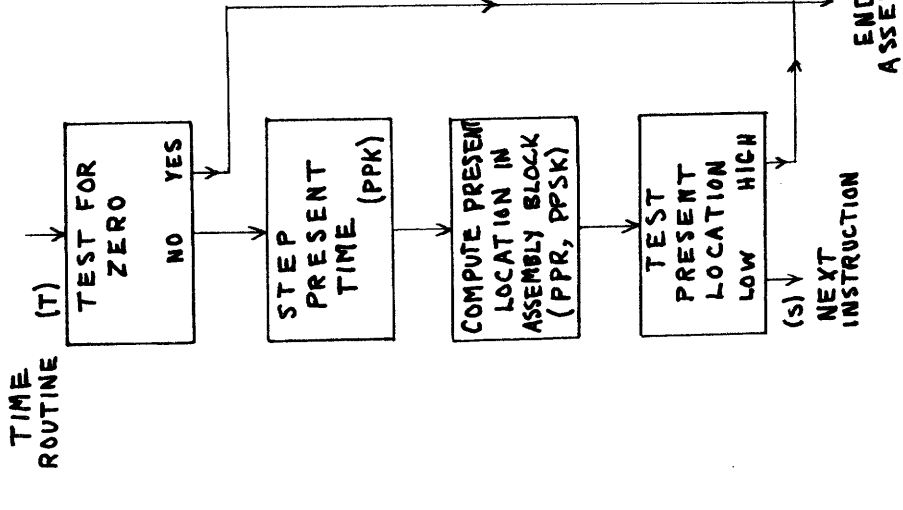
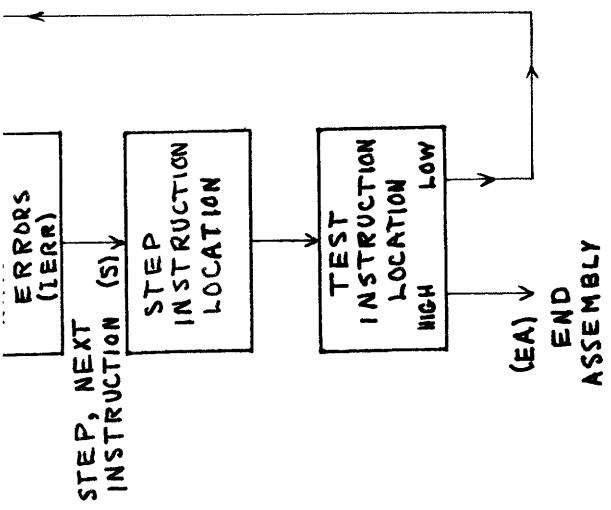


NOISE ROUTINE



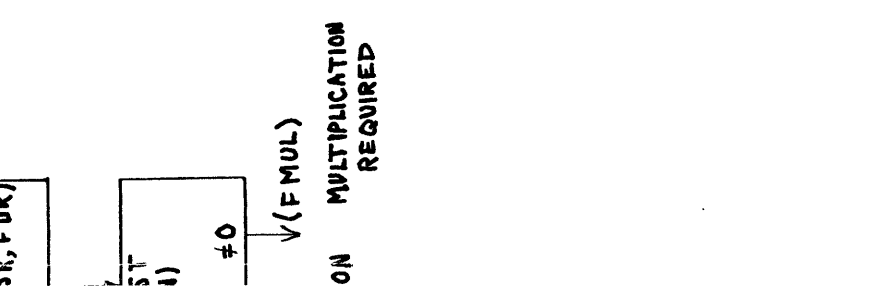
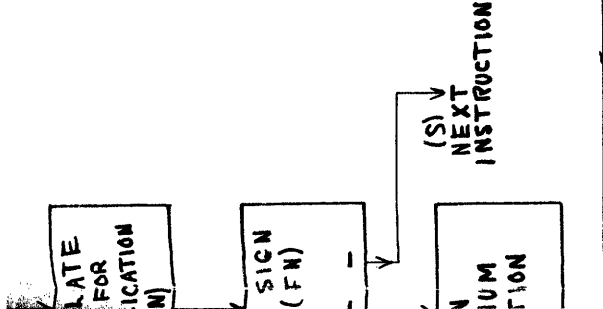
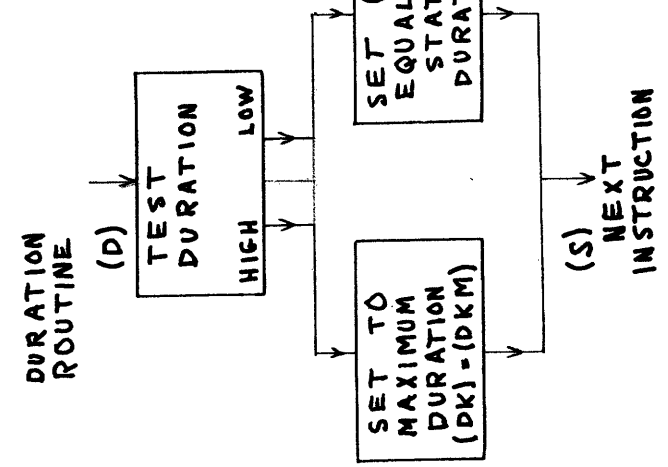
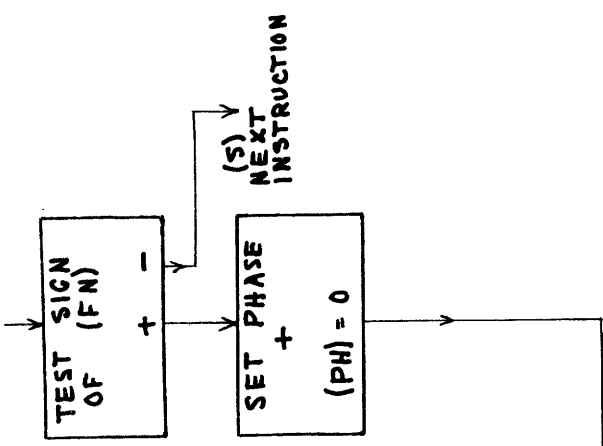
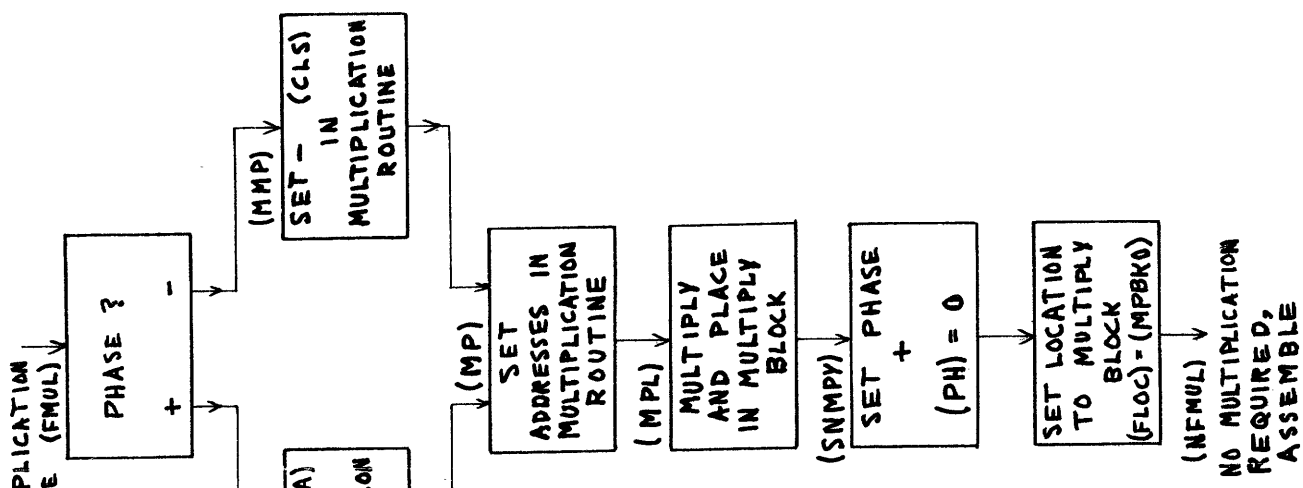
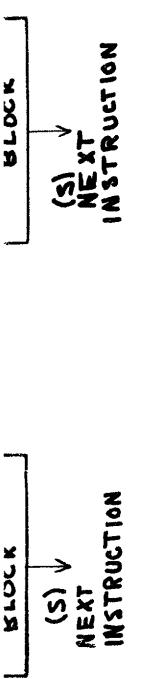
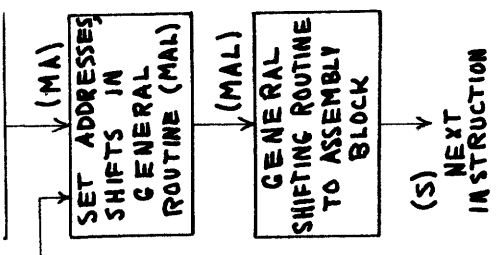
NO MULTIPLICATION REQUIRED (NFMUL)



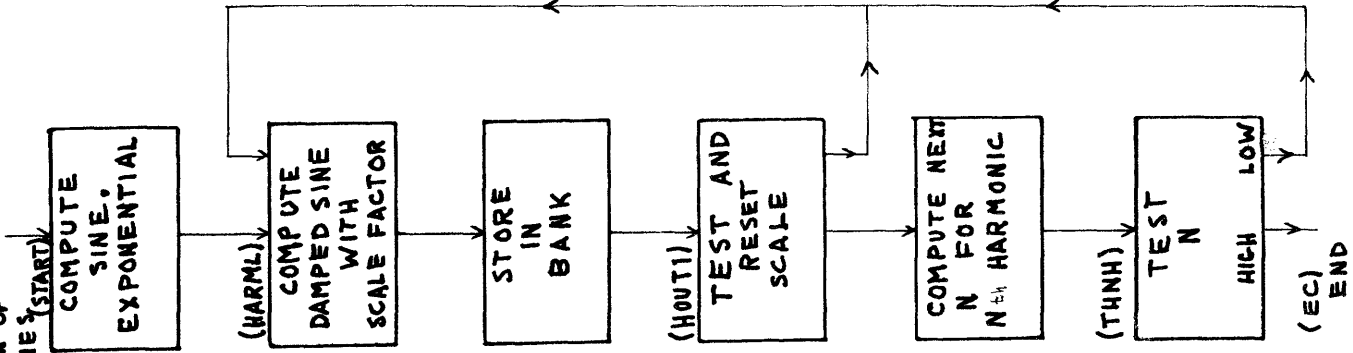


GENERAL SYNTHESIS PROGRAM

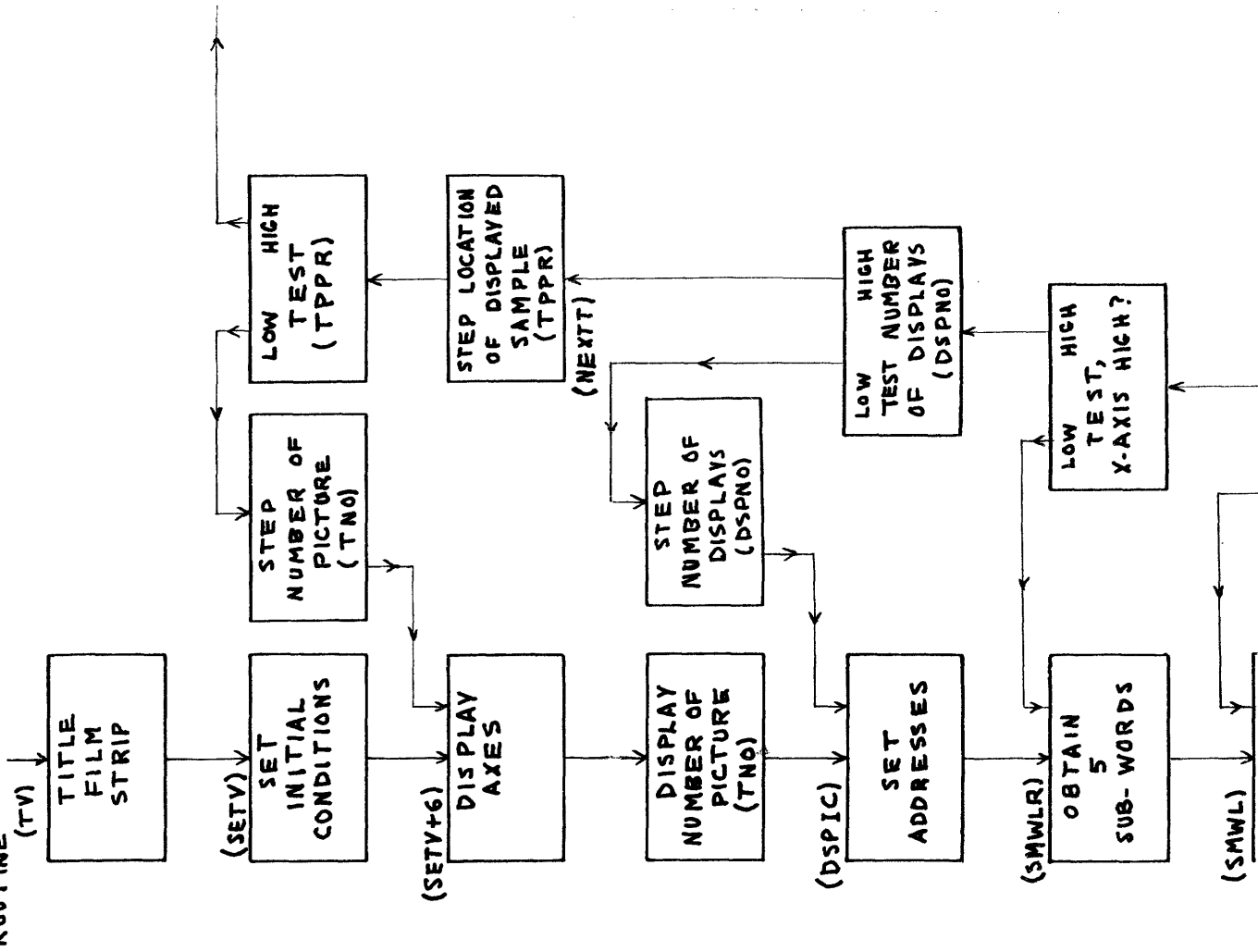
SYMBOLS IN PARENTHESES REFER TO LOCATIONS IN THE SAP PROGRAM



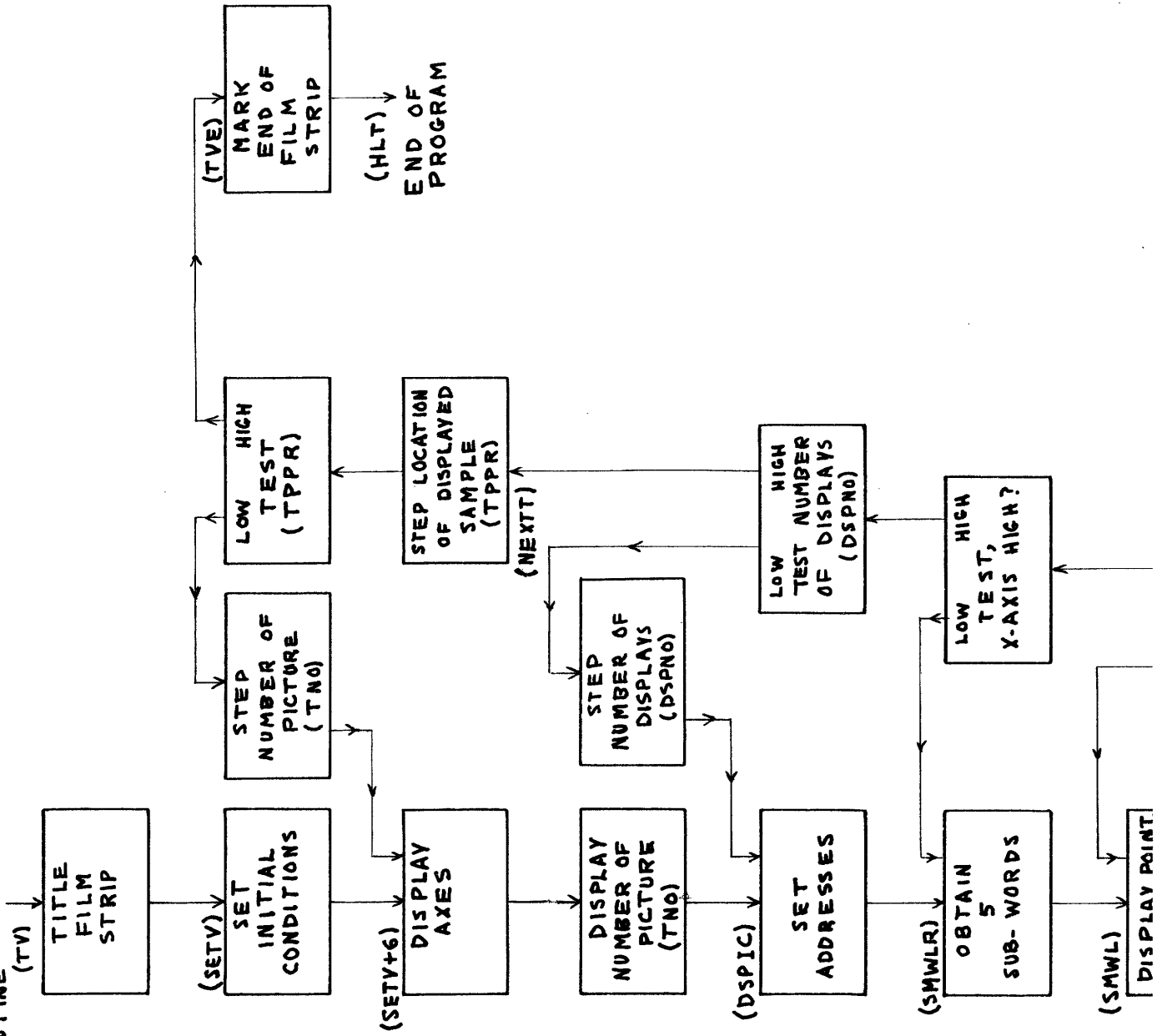
INITIAL
COMPUTATION OF
DAMPED SINE



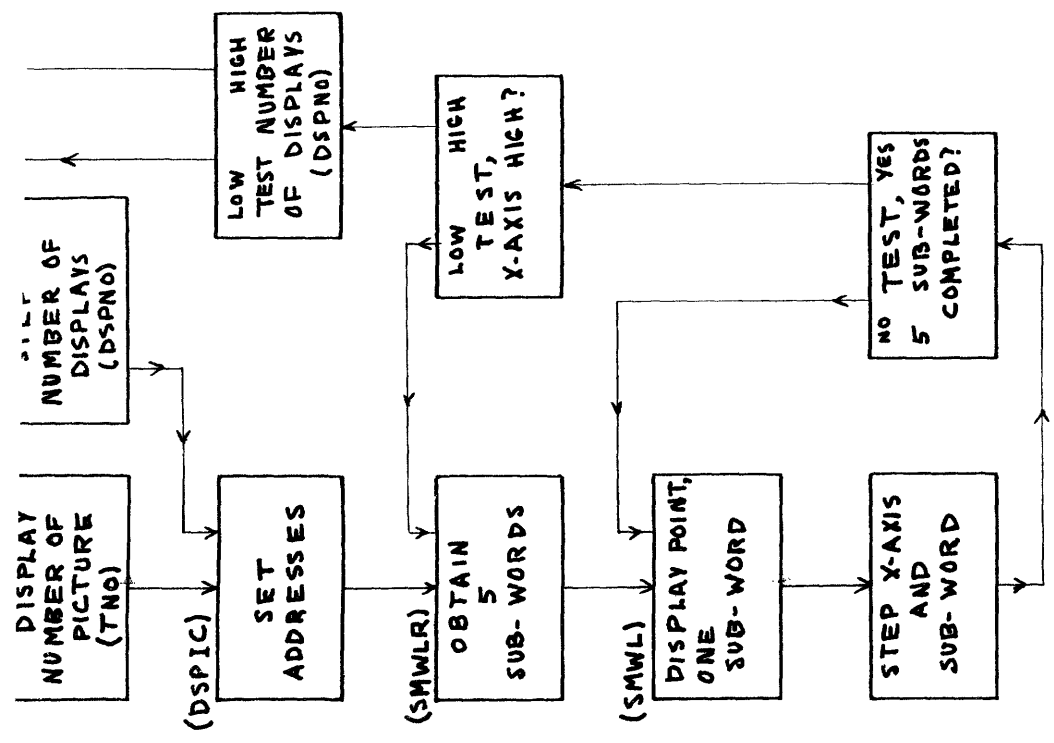
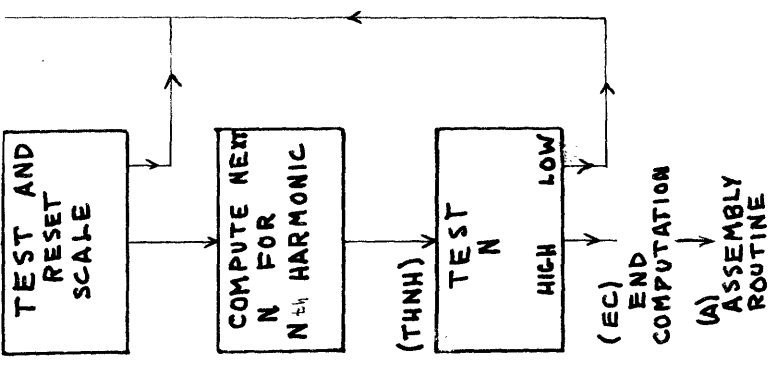
TV
DISPLAY
ROUTINE



TV
DISPLAY
ROUTINE

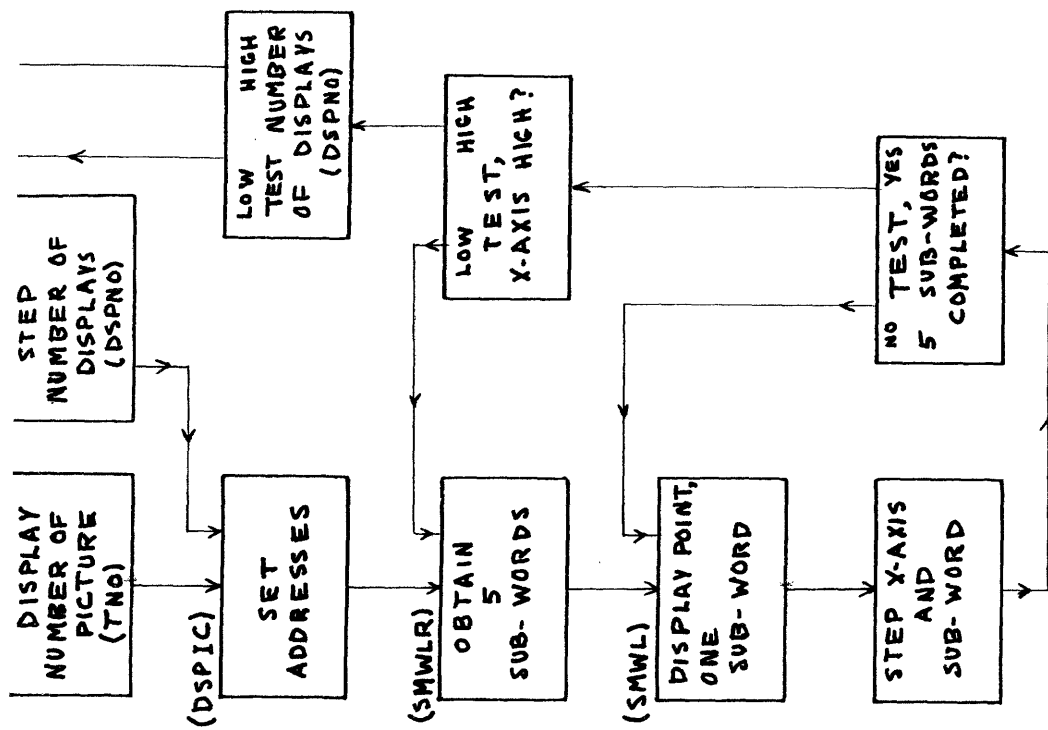


AUXILIARY ROUTINES



AUXILIARY ROUTINES

SYMBOLS IN PARENTHESES
REFER TO LOCATIONS IN
THE SAP PROGRAM



Appendix D

SAP - PROGRAM

Index of the more important entry points.

*

General Synthesis Program

<u>Symbol</u>	<u>Instruction Number</u>
A	80
D	148
E	103
EA	106+1
F	155
FMUL	210
N	385
NFMUL	202
MA	297
MAL	324
MDA	244
MDAL	263
MMA	279
MMAL	289
MMP	346
MP	352
MPL	368
PDA	213
PDAL	232

Program Index (continued)

<u>Symbol</u>	<u>Instruction Number</u>
PMP	349
PSKO	207
S	107
SNMPY	381
T	133
TD	193

Auxiliary Routines

A	80
DSPIC	427
EC	79
HARML	27
HOUT1	54
NEXTT	458
SETV	413
SMWL	434
SMWLR	430A
START	1
THNH	75
TV	405
TVE	468

THE SYNTHESIS PROGRAM

AS RUN, MAY, 1959

START	CAL	SINOL		1
	PAX	0,1		2
	PXD			3
	TRA	**4		4
	CAL	BETA		5
	ADM	DBETA		6
	SLW	BETA		7
	TSX	SIN2,4		8
	STO	SINO+256,1		9
	TIX	DSINO,1,1		10
	CLA	EXPOL		11
	PAX	0,1		12
	PXD			13
	TRA	**4		14
	CLA	EXPF		
	FAD	DEXP		
	STO	FXPF		
	TSX	S816EX,4		
	NOP			
	UFA	FIX		
	ROL	9		
	LGL	8		
	STO	EXP0+320,1		20
	TIX	DEXP0,1,1		21
	CAL	HN		22
HNCOM	COM			23
	ADD	I1		24

LOOP

EXPONENTIAL COMPUTATION

LOOP

ALS 18	25
STD HTN	26
HARML LXA HIR10,1	27
LXA HIR20,2	28
LXA HIR40,4	29
HAR1 LDQ SIN0+255,1	30
MPY FXP0+320,2	31
LRS 35	32
SCALE MPR HSCAL+2	32A
ADM 1CRT	33
ANA 1CMASK	34
STO 5WRD+5,4	35
TNX H5WRD,4,1	36
HTN TXI *+1,1,-1	37
TXI *+1,2,-1	38
PXD 0,1	39
ARS 18	40
ANA MDMSK	41
PAX 0,1	42
TRA HAR1	43
H5WRD CAL 5WRD+5,4	44
LRS 7	45
TXI *+1,4,1	46
TXL *-3,4,5	47
CAL SQ5WD	48
ADM 11	49
STA SQ5WD	50
SQ5WD STO FBLK-1	51
TXL HOUT1,2,1	52
TRA HAR1	53
HOUT1 CLA TSC	54
TMI SHN	55
SSM	56
STO TSC	57
CLA HSCAL+1	58
STA SCALE	59
TRA HARML	60

COMPUTE SINE TIMES EXPONENTIAL

SCALE TO DESIRED AMPLITUDE

STEP N, MODULO SINE LENGTH

FORM DAMPED SINES IN STORAGE

RESET, SCALE 2

3

SHN	SSP			61
	STO	TSC		62
	CLA	HSCAL		63
	STA	SCALE		64
	CLA	HN	RFSFT, SCALF 1	66
	CAS	HN1		67
	NOP			68
SHN1	TRA	SHN2	HN GRTR THAN HN1	69
	ADD	I1	HN LESS THAN HN1	70
	STO	HN		71
	TRA	THNH		72
SHN2	ADD	I10	INITIAL TEST VALUE	73
	STO	HN		74
	CAS	HN2		75
	TRA	EC		76
	NOP			77
	TRA	HNCOM		78
EC	TRA	A		79
A	LDG	PPR0	FND COMPUTATION OF SINE, EXP.	80
	STG	PPR	SET PPR, ASSEMBLY INITIAL CONDITIONS	81
	MPY	I5	5*PPR	82
	STG	PPK		83
	STZ	PP5K		84
	STZ	ILT		85
	CLA	INS00	INITIALIZES INS BLOCK	86
	STA	S		87
	TRA	S		88
TINS	LDG	INS		89
	PXD			90
	LGL	3		91
	ADD	TINS0		92
	STA	*+1		93
	TRA	*		94
TRINS	TRA	T	TRANSFER TO PROPER INSTRUCTION ROUTINE	95
	TRA	D		96
	TRA	F		97
	TRA	N		98

TRA E	99
TRA F	100
TRA E	101
TRA E	102
CLA IFERR	103
ADD I1	104
STO IFERR	105
TRA S	106
EA	
TRA TV	
S	
CLA INSO	107
STO INS	108
CLA ILT	109
TMI INSM	110
SSM	111
STO ILT	112
CAL INS	113
LRS 18	114
LDO IO	115
LRS 18	116
LGL 1	117
STQ INS	118
TRA IMLT	119
INSM	120
SSP	121
STO ILT	122
LDO INS	123
LGL 18	124
STQ INS	125
CAL S	126
ANA MSKA	126A
ADD I1	126B
STA S	127
IMLT	128
CAL S	129
ANA MSKA	130
CAS INSOM	131
NOP	132
TRA EA	
TRA TINS	

INS ERROR

END ASSEMBLY, TRA TO TV
OBTAIN INSTRUCTION (INS)

LEFT HALF INS

RIGHT HALF INS

TEST MAX INS LOC

END ASSEMBLY

T	PXD	PRESENT TIME ROUTINE	133
	LGL 15		134
	TZF FA	T=0, END ASSEMBLY	135
	ADD PPK		136
	STO PPK	RESET PPK	137
	LRS 35		138
	DVP I5		139
	STO PPR	FORM PPR	140
	STO PPSK	FORM PPSK	141
	PXD		142
	LLS 35		143
	CAS PPRM	TEST FOR END OF ASSEMBLY BLOCK	144
	NOP		145
	TRA EA	END ASSEMBLY	146
	TRA S		147
	PXD	DURATION ROUTINE	148
	LGL 15		149
	CAS DKM	TEST FOR D MAXIMUM	150
	NOP		151
	CLA DKM		152
	STA DK	SET DK	153
	TRA S		154
	LGL 2	FORMANT ROUTINE	155
	PXD		156
	LLS 0		157
	STO PH	SET PHASE	158
	LGL 1		159
	PXD 6		160
	CAS FM		161
	NOP		162
	TRA E	FREQ. HIGH	163
	ALS 1		164
	STO COMMON		165
	PXD		166
	LGL 6	STORE AMPLITUDE	167
	STO COMMON+1		168
			169

35

LRS 1		170
PXD		171
LLS 1	COMPUTF FLOC	172
ADD COMMON		173
LRS 35		174
MPY FWDL		175
LLS 35		176
ADD FO		177
STA FLOC		178
CLA COMMON+1		179
ARS 1		180
ANA I7		181
ADD FNO		182
STA *+1		183
CLA *	OBTAIN N FOR MPY	184
STO FN		185
TMI S		186
CLA FDKMA		187
ADD FN		188
STA *+1		189
CLA *	OBTAIN DURATION	190
CAS DK		191
NOP		192
CLA DK		193
STO FDK	SET MAXIMUM DURATION	194
LRS 35		195
DVP I5		196
STO FDR	SET FDR	197
STO FDSK	SET FDSK	198
CLA FN		199
TZE NFMUL	NO MULTIPLY	200
TRA FMUL	MULTIPLY	201
CLA PPSK	NO MULTIPLY	202
TZE PSKO	PPSK=0	203
CLA PH		204
TPL MA	PLUS TO ASSEMBLY	205
TPA MMA	MINUS TO MPBLK	206

PSK0	CLA PH			207
	TPL PDA	PLUS DIRECT ASSEMBLY		208
	TRA MDA	MINUS DIRECT ASSEMBLY		209
FMUL	CLA PH			210
	TMI MMP	MINUS MULTIPLY		211
	TRA PMP	PLUS MULTIPLY		212
PDA	CLA FDR	PPSK=0, PLUS DIRECT ASSEMBLY, SET-UP		213
	PAX 0,1			214
	ADM FLOC			215
	STA PDAL	SET B		216
	STA PDAL+5			217
	CLA FDR			218
	ADM PPR			219
	STA PDAL+1	SFT A		220
	STA PDAL+2			221
	STA PDAL+9			222
	STA PDAL+10			223
	CLA I5	COMPUTE (5-FDSK)*7		224
	SUR FDSK			225
	LRS 35			226
	MPY I7			227
	LLS 35			228
	STA PDAL+6			229
	STA PDAL+8			230
	TRA PDAL			231
PDAL	CLA *,1	R, PLUS DIRECT ASSEMBLY LOOP		232
	ADM *,1	A		233
	STO *,1	A		234
	TIX PDAL,1,1			235
	TXI *,1,1,-1			236
	CLA *,1	B		237
	LRS 0	(5-DSK)*7		238
	LDO IO			239
	LLS 0	(5-DSK)*7		240
	ADM *,1	A		241
	STO *,1	A		242
	TRA S			243

MDA	CLA FDR	PPSK=0, MINUS DIRECT ASSEMBLY, SET-UP	244
	PAX 0,1		245
	ADM FLOC		246
	STA MDAL	SET B	247
	STA MDAL+7		248
	CLA FDR		249
	ADM PPR		250
	STA MDAL+3	SET A	251
	STA MDAL+4		252
	STA MDAL+13		253
	STA MDAL+14		254
	CLA I5	COMPUTE (5-FDSK)*7	255
	SUB FDSK		256
	LRS 35		257
	MPY I7		258
	LLS 35		259
	STA MDAL+8		260
	STA MDAL+10		261
	TRA MDAL		262
MDAL	CLS *,1	R, MINUS DIRECT ASSEMBLY LOOP	263
	ADD CRITS		264
	ANA CRITO		265
	ADM *,1	A	266
	STO *,1	A	267
	TIX MDAL,1,1		268
	TXI *,1,1,-1		269
	CLA *,1	B	270
	LRS 0	(5-DSK)*7	271
	LDQ MIO		272
	LLS 0	(5-DSK)*7	273
	ADD CRITS		274
	ANA CRITO		275
	ADM *,1	A	276
	STO *,1	A	277
	TRA S		278
MMA	CLA FDR	MINUS TO MPBK, SET-UP	279
	ADM I1		280

38

PAX 0,1			281
ADM FLOC			282
STA MMAL	SET B		283
CLA FDR			284
ADM MPRK0			286
STA MMAL+3	SET C		287
TRA MMAL			288
CLS *,1	B, OBTAIN F, MINUS TO MPBK		289
ADD CRITS	FORM NEGATIVE		290
ANA CRITO			291
STO *,1	C		292
TIX MMAL,1,1			293
CLA MPRK0	RESET FLOC		294
STA FLOC			295
TRA MA			296
CLA FDR	SFT-UP FOR MAL		297
PAX 0,1			298
ADM FLOC			299
STA MAL	SET B		300
STA MAL+9			301
CLA FDR			302
ADM PPR			303
STA MAL+2	SET A		304
STA MAL+3			305
STA MAL+15			306
STA MAL+16			307
STA MAL+19			308
STA MAL+20			309
LDO PPSK			310
MPY I7			311
LLS 35			312
STA MAL+4	SET PPSK*7		313
SUB I35			314
STA MAL+1	-(5-PPSK)*7		315
STA MAL+14			316
LDO FDSK			317
MPY I7			318

LLS 35						319
STA MAL+10				SET FDSK*7		320
STA MAL+12						321
TRA MAL-1				END MAL SET-UP		322
PXD						323
MAL	LDQ *,1	R, GENERAL	ASSEMBLY SHIFTING LOOP			324
	LLS 0		(5-PPSK)*7			325
	ADM *,1	A				326
	STO *,1	A				327
	LLS 0		PPSK*7			328
	TIX MAL,1,1					329
	STO COMMON					330
	PXD					331
	TXI *,1,1,-1					332
	LDQ *,1	B				333
	LLS 0		DSK*7			334
	LDQ IO					335
	LRS 0		DSK*7			336
	CLA COMMON					337
	LLS 0		(5-PPSK)*7			338
	ADM *,1	A				339
	STO *,1	A				340
	LLS 35					341
	TXI *,1,1,-1					342
	ADM *,1	A				343
	STO *,1	A				344
	TRA S		END MAL			345
MMP	CLA CLS		SET MINUS MULTIPLY			346
	STO MPL					347
	TRA MP					348
PMP	CLA CLA		SET PLUS MULTIPLY			349
	STO MPL					350
	TRA MP					351
MP	CLA FDR		MULTIPLY ROUTINE, SET-UP			352
	ADD I1					353
	PAX 0,1					354
	ADD FLOC					355

STA MPL	SET FOR FLOC	356
SUR FLOC		357
ADD MPRKO		358
STA MPL+10	SET FOR MPB LOCATION	359
CLA FN	SET SHIFT, N	360
STA MPL+2		361
ADD JNKNO		362
STA MPL+3	SET JUNK MASK (N)	363
CLA FN		364
ADD MSKNO	SET SIGN MASK (N)	365
STA MPL+5		366
TRA MPL		367
MPL	MULTIPLY ROUTINE	368
CLS *,1		369
ADD CRITS		370
ARS 0	N	371
ANA *	JUNKN	372
STO COMMON		373
ANA *		374
ALS 1		375
SUR CBITS		376
SUR COMMON		377
ANA CRITO		378
SLW *,1	MPRK	379
TIX MPL,1,1		380
TRA SNMPY		381
SNMPY STZ PH	SET FOR NO MULTIPLY, SET PHASE +	382
CLA MPRKO		383
STO FLOC	SET FLOC	384
TRA NFMUL		385
PXD 0	NOISE ROUTINE	386
LGL 6		387
CAS NM		388
NOP		389
TRA E	FREQUENCY HIGH, ERROR	390
STO COMMON		391
PXD 0		392
LGL 3		

ADD NNO	393
STA *+1	394
CLA *	395
TMI S	396
STO FN	397
LDQ COMMON	398
MPY NWDL	399
LLS 35	400
ADD NO	401
STA FLOC	402
STZ PH	403
TRA TD	404
WTV	405
CFF	406
TSX WRIMG,4	407
PZE 500,7,100	408
PZE BCDT,0,5	409
STZ BCDTE	410
CFF	411
TRA SFTV	412
SETV	413
CLA PPRO	414
ADD I100	415
STA SMWL-2	416
STA TPPR	417
CLA I1	418
STO TNO	419
CPY XAXIS	420
CPY YAXIS	421
CPY XMARK	422
LDQ TNO	423
TSX WRINT,4	424
PZE 770,7,770	425
STZ TVERR	426
STZ DSPNO	427
DSPIC CLA I100	428
PAX 0,1	429
PXD 0	

SET FN
COMPUTE ADDRESS

SET ADDRESS

TV DISPLAY ROUTINE

TITLE TV RUN

ERROR

INITIALIZF BLOCK ADDRESS

SET TNO=1
X-Axis
Y-Axis
RIGHT EDGE MARK

NUMBER PHOTO

ERROR
SET DIPLAY NO. = 0
DISPLAY, INITIAL CONDITIONS
LARGE WORDS

PAX 0,6	SMALL WORDS AND X-AXIS	430
SMWLR PXD 0		430A
PAX 0,2		430B
TRA SMWL-2		431
LDQ ARLK+100,1	ORTAIN LARGE WORD	432
LGL 1		433
SMWL LGL 1	SMALL WORD LOOP	434
LLS 0		435
CHS		436
LRS 0		437
PXD 0		438
LGL 6	SMALL WORD IN AC	439
ALS 4	SCALE	440
STA TVPT		441
SXD TVPT,4		442
STQ COMMON	SAVE MQ FROM CPY DISTURBANCE	443
CPY TVPT	DISPLAY POINT	444
NOP	DISPLAY DELAY (TEMPORARY)	
NOP		
NOP		
NOP		
NOP		
NOP		
NOP		
LDQ COMMON		445
TXI *+1,4,2	STEP X-AXIS	446
TXI *+1,2,1	STEP SMALL WORD	447
TXH *+2,2,4		448
TRA SMWL		448A
TIX SMWLR,1,1	NEXT LARGE WORD	449
CLA DSPNO		450
CAS DSPNM		451
NOP		452
TRA NEXTT	NO. OF DISPLAYS HIGH, NEW PICTURE	454
ADD I1		455
STO DSPNO	STEP DISPLAY NO.	456
TRA DSPIC	NEXT DISPLAY	457

NEW PICTURE

NEXTT CLA TPRR 458
ADD I100 459
STA SMWL-2 460
STA TPRR 461
CAS PPRM 462
NOP 463
TRA TVE1 464
CLA TNO 465
ADD I1 466
STO TNO 466A
TVCAM NOP 466B
TRA SFTV+6 467
TVE1 HPR 77 467A
SWT 1 B
TRA TVE C
CAL CFF D
SLW TVCAM E
TRA TV F

TEST FOR FILM RECORDING

NO, END
YES, INSERT CFF

TITLE END OF TV RUN

TVE CFF 468
TSX WRIMG,4 469
PZF 500,7,100 470
PZF RCDTF,0,4 471
STZ BCTFE 472
CFF 473
HLT 474
RFM D1
PANIC * * * STOP * * * END OF PROGRAM D1A
DATA FOLLOWS

DSPNM DEC 10
DSPNO DEC 0
TNO DEC 0
TPPR DEC 0
FDSK DEC 0
FLOC DEC 0
PPSK DEC 0
I0 DEC 0
I1 DEC 1
I2 DEC 2
I3 DEC 3

D2
D3
D4
D5

I4	DEC 4	D6
I5	DEC 5	D7
I6	DEC 6	D8
I7	DEC 7	D9
I10	DEC 10	D10
I35	DEC 35	D11
I100	DEC 100	D11A
I200	DEC 200	D12
M10	DEC -0	D13
IFRR	DFC 0	D14
TVERR	PZF -1	D15
BCDTE	PZE -1	D16
BCTFE	PZE -1	D18
REM		D19
BETA	DEC 0	D20
EXPF	DEC 0	D21
DBETA	DFC 1R8	D22
DEXP	DEC -1.08E-2	D23
SIN0L	DEC 256	D24
EXP0L	DEC 320	D25
FIX	OCT 2010000000000	D26
HN	DEC 2	D27
ICRT	OCT 100	D28
ICMSK	OCT 77	D29
MDMSK	OCT 377	D30
HIR10	DEC 255	D31
HIR20	DEC 320	D32
HIR40	DEC 5	D33
HSCAL	PZE #+2	D34
	PZE #+2	D35
	DEC 40	D36
	DEC 60	D37
TSC	DEC 0	D38
HN1	DFC 10	D39
HN2	DEC 105	D40
REM		D41
TINSO	PZE TRINS	D42

ERRORS

SINE AND EXP. COMPUTATION

INITIAL TEST VALUE

ASSEMBLY

45

ILT	DEC 0	D44
INS	DEC 0	D45
MSKA	PZE -1	D46
INSOM	PZE INS0+2000	D47
INS00	PZE INS0	D48
PPK	DEC 0	D49
PPR	DEC 0	D50
PPR0	PZE ARLK	D51
PPRM	PZE ARLK+2000	D52
DK	DEC 0	D53
DKM	DEC 1280	D54
PH	DEC 0	D55
FM	DEC 22	D56
FWDL	DEC 64	D57
FO	PZE FRLK	D58
FN	DEC 0	D59
FDKM	DEC 320	D60
	DEC 255	D61
	DEC 190	D62
	DEC 125	D63
	DEC 65	D64
FDKMA	PZE FDKM	D65
FNO	PZE *+1	D67
	DEC -0	D68
	DEC 4	D69
	DEC 3	D70
	DEC 2	D71
	DEC 1	D72
	DEC 0	D73
	DEC 0	D74
	DEC 0	D75
FDK	DEC 0	D76
FDR	DEC 0	D77
CRITS	OCT 201004020100	D78
CRITO	OCT 176773757677	D79
MPRK0	PZE MPRK	D80
NNO	PZE *+1	D81

INITIAL TEST VALUE

INITIAL TEST VALUE

DEC -0						D82	
DEC 5						D83	
DEC 4						D84	
DEC 3						D85	
DEC 2						D86	
DEC 1						D87	
DEC 0						D88	
DEC 0						D89	
DEC 30						D90	
PZE NBLK						D91	
DEC 256						D92	
DEC 1280						D93	
CLA #,1						D94	
CLS #,1						D95	
CFF						D95A	
RFM						D96	
BCDT	BCD	5SSI	1	R.P.	FUTRELLE	M630-2-1	D97
BCDTF	BCD	4SSI	1	M630-2-1	END		D98
XAXIS	MTW	512					D99
YAXIS	MON	0					D100
TVPT	DEC	0					D100A
XMARK	MON	0,0,1000					D101
RFM							D102
JNKNO	PZE	*					D103
OCT	376371747637						D104
OCT	076371747637						D105
OCT	036170743617						D106
OCT	016070341607						D107
OCT	006030140603						D108
OCT	002010040201						D109
MSKNO	PZE	*					D110
OCT	040201004020						D111
OCT	020100402010						D112
OCT	010040201004						D113
OCT	004020100402						D114
OCT	002010040201						D115
RFM							D116

INITIAL TEST VALUE

TV

MASKS

LIBRARY ROUTINES

47

NA.110	LIB								D117
NA.117	LIB								D118
LAS816	LIB								D119
SIN2	RFM MU SIN 2	0039	CARDS	FIXED POINT	SINE				MUSIN201
	SXD COMMON,1			SAVE	IR1				SIN20002
	ALS 1			THETA/PI					SIN20003
	PBT 1			TEST FOR THETA/2PI	GREATER OR =1/2				SIN20004
	TRA SIN2+5								SIN20005
	CHS								SIN20006
	ALS 1								SIN20007
	PBT 1			TEST FOR THETA/PI	GREATER OR =1/2				SIN20008
	TRA SIN2+9								SIN20009
	COM								SIN20010
	LRS 1								SIN20011
	RND								SIN20012
	STO COMMON+1								SIN20013
	LDQ COMMON+1			THETA/PI					SIN20014
	MPY COMMON+1			(THETA/PI)EXP2					SIN20015
	LLS 2			(2THETA/PI)EXP2					SIN20016
	RND								SIN20017
	PBT 1			TEST FOR OVERFLOW					SIN20018
	TRA SIN2+19								SIN20019
	COM								SIN20020
	STO COMMON+2								SIN20021
	LXA SIN2+18,1			SET	IR1=6				SIN20022
	LDQ SIN2+31			A	SUR 13				SIN20023
	MPR COMMON+2			(2 THETA/PI)EXP2					SIN20024
	ADD SIN2+38,1								SIN20025
	LRS 35								SIN20026
	TIX SIN2+22,1,1			LOOP TRANSFER, OUT	AFTER 6 PASSES				SIN20027
	MPY COMMON+1			THETA/PI					SIN20028
	LLS 1								SIN20029
	RND								SIN20030
	LXD COMMON,1			RESTORE	IR1				SIN20031
	TRA 1,4			OUT					SIN20032
	OCT 000000001650			A	SUB 13				SIN20033
	OCT 400000170505			A	SUB 11				SIN20034

OCT 000012407327	A SUB 9	SIN20035
OCT 400462645440	A SUB 7	SIN20036
OCT 012146570632	A SUB 5	SIN20037
OCT 522535716304	A SUB 3	SIN20038
OCT 311037552421	A SUB 1	SIN20039

BLOCKS OF STORAGE

COMMON	BSS 20
5WRD	BSS 20
MPBK	BSS 400
SINO	BSS 300
FXPO	BSS 350
FBLK	BSS 7000
ARLK	BSS 7000
NRLK	SYN FRLK
	FND START

ONLY INITIALLY

Appendix E

The Format of the Instructions

Octal numbers are based on the numbers 0 through 7 as our decimal system is based on 0 through 9. Octal numbers are especially useful in computer work as there is a one-to-one correspondence between an octal digit and a group of three binary digits. The instructions will be composed of six octal digits so that two will occupy each 704 word. The positions are numbered from the left.

- T will have a 0 in position 1. Positions 2-6 contain the value of the shift in units of one sub-word, from 0 to $77,777_8$ (0 to $32,767_{10}$.) A 0 will terminate assembly.
- D will have a 1 in position 1. Positions 2-6 contain the duration in units of one sub-word.
- F will have a 2 in position 1. Position 2 is even for a starting phase of 0° (plus) and odd for a starting phase of 180° (minus.) Positions 3 and 4 contain the number of the sample and thus determines the frequency. Positions 5 and 6 contain the amplitude.
- N will have a 3 in position 1. Positions 2 and 3 specify the nature of the noise desired. The samples would be numbered and their characteristics listed for the programmer's convenience. Position 4 specifies the amplitude. Positions 5 and 6 are disregarded.

Appendix F - Instructions

THE INSTRUCTIONS USED TO TEST

THE SYNTHESIS PROGRAM

INSO	OCT 000062100062	T50, D50
	OCT 201212000144	F+550, 30DB, T100
	OCT 211212000144	F-550, 30DB, T100
	OCT 201210000144	F+550, 24DB, T100
	OCT 211210000310	F-550, 24DB, T200
	OCT 100024201212	D20, F+550, 30DB
	OCT 000050100027	T40, D23
	OCT 201212000050	F+550, 30DB, T40
	OCT 211212000050	F-550, 30DB, T40
	OCT 000001100024	T1, D20
	OCT 201212000050	F+550, 30DR, T40
	OCT 211212000050	F-550, 30DR, T40
	OCT 100027201212	D23, F+550, 30DR
	OCT 000050211212	T40, F-550, 30DR
	OCT 000403110000	T259, DMAX
	OCT 304600001440	N4(550,1100) 36DB, T800
	OCT 100226306500	D150, N6(2750,3600) 30DB
	OCT 000062100702	T50, D450
	OCT 200212201212	F+110, 30DB, F+550, 30DR
	OCT 000764200210	T500, F+110, 24DB
	OCT 201210000764	F+550, 24DB, T500
	OCT 211212200212	F-550, 30DR, F+110, 30DR
	OCT 000764211212	T500, F-550, 30DR
	OCT 210212201312	F-110, 30DB, F+1100, 30DR
	OCT 000764100024	T500, D20
	OCT 201312201512	F1100, 30DR, F2200, 30DR
	OCT 100024000036	D20, T30

OCT 201312201512	F1100, 30DB,	F2200, 30DB
OCT 100024000036	D20, T30	
OCT 201312201512	F1100, 30DB,	F2200, 30DB
OCT 100024000036	D20, T30	
OCT 201312201512	F1100, 30DB,	F2200, 30DB
OCT 100024000036	D20, T30	
OCT 201312201512	F1100, 30DB,	F2200, 30DB
OCT 100024000036	D20, T30	
OCT 201312201512	F1100, 30DB,	F2200, 30DB
OCT 100024000036	D20, T30	
OCT 201312201512	F1100, 30DB,	F2200, 30DB
OCT 100024000036	D20, T30	
OCT 201312201512	F1100, 30DB,	F2200, 30DB
OCT 100024000036	D20, T30	
OCT 201312201512	F1100, 30DB,	F2200, 30DB
OCT 100024000036	D20, T30	
OCT 201312201512	F1100, 30DB,	F2200, 30DB
OCT 0	THE VERY END	