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FENCHEL AND LAGRANGE DUALITY ARE EQUIVALENT

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ABSTRACT

A basic result in ordinary (Lagrange) convex programming is the saddlepoint duality theorem concerning optimization problems with convex inequalities and linear-affine equalities satisfying a Slater condition. This note shows that this result is equivalent to the duality theorem of Fenchel.

Among the most powerful tools in mathematical programming are the saddlepoint duality theories for both Fenchel and ordinary (Langrange) convex programs. The purpose of this note is to exhibit the equivalence of these two theories by showing that each can be used to develop the other.

The connection between Fenchel and Lagrange duality is certainly not surprising since each is based upon separating hyperplane arguments. Previously, Whinston [7] has shown how the Lagrangian theory can be derived from Frenchel's results when optimizing over Rⁿ in the absense of equality constraints. Both theories also can be obtained from Rockafellar's general perturbation theory. The full equivalence between these theories, though, has not been stated in the open literature and seems to be unknown to most of the mathematical programming community. In fact, a number of recent comprehensive books in convex optimization including Luenberger [3], Rockafellar [5] and Stoer and Witzgall [6] do not explicitly mention or exploit this equivalence, but rather develop the two theories separately.

The Fenchel and Lagrange convex programs can be stated as

Each C_j is a convex subset of R^n ; $f_0(x)$ and $f_1(x)$ are convex functions defined respectively on C_0 and C_1 ; $f_2(x)$ is a concave function defined on C_2 ; $g:C_0 \to R^m$ is convex and $h:R^n \to R^r$ is linear-affine.

The saddlepoint dual problems are:

Fenchel Dual

d=max {
$$f_2^*(\pi) - f_1^*(\pi)$$
}

 $\pi \in \mathbb{R}^n$

Lagrange Dual

 $d=\max \inf_{\pi \geq 0} \{L(x;\pi,\alpha)\}$
 $\pi \geq 0$
 $\alpha \in \mathbb{R}^r$
 $\pi \in \mathbb{R}^n$

where

$$f_1^*(\pi) = -\inf \{f_1(x) - \pi x\}, f_2^*(\pi) = \inf \{\pi x - f_2(x)\}$$
 $x \in C_1$

are respectively conjugate convex and conjugate functions and

$$L(x;\pi,\alpha) = f(x) + \pi g(x) + \alpha h(x)$$

is the Lagrangian function. Note that the dual problems are formulated as maximizations. We shall consider conditions to insure that these maxima are attained.

Recall that $f_1^*(\pi) = +\infty$ and/or $f_2^*(\pi) = -\infty$ are possibilities. Such values for π will be inoperative for the maximization in the Fenchel dual and thus that problem is frequently written with $\pi \epsilon C_1^* \cap C_2^*$ instead of $\pi \epsilon R^n$ where

$$C_1^* = \{ \pi \in \mathbb{R}^n : f_1^*(\pi) < +\infty \} \text{ and } C_2^* = \{ \pi \in \mathbb{R}^n : f_2^*(\pi) > -\infty \}.$$

Two of the most useful and delicate theorems connecting these primal-dual pairs are (for any convex C, ri(C) denotes its relative interior):

Theorem 1: If $v > -\infty$ and $ri(C_1) \cap ri(C_2) \neq \emptyset$, then v=d.

Theorem 2: Assume $\overline{v} > -\infty$ and the Slater condition that there exists an $\mathbf{x}^0 \in \mathrm{cri}(C_0)$ such that $\mathrm{g}(\mathbf{x}^0) < \mathrm{o}$ and $\mathrm{h}(\mathbf{x}^0) = \mathrm{o}$. Then $\overline{v} = \overline{\mathrm{d}}$

The first theorem is Fenchel's original duality theorem [2] and the second is a consequence of a result due to Fan, Glicksberg and Hoffman [1] (see Chapter 28 of [5] or Chapter 6 of [6]). When there are no inequality constraints, theorem 2 remains valid by omitting the reference to g(x) in the Slater condition. $\overline{v}=\overline{d}$ is equivalent to saying that the Kuhn-Tucker condition is satisfied: there exists $\hat{\pi}>0$ and $\hat{\alpha}$ such that $\overline{v}\leq L(x,\hat{\pi},\hat{\alpha})$ for all $x\in C_0$.

Below we show that either theorem 1 or 2 can be used to derive the other. The following elementary result concerning relative interiors will be useful.

Proof: The following condition ([5] theorem 6.4) characterizes the relative interior of any convex set S:

zeri(S) if and only if for each $\bar{z} \in S$ there is a $\mu>1$ such that $\mu z+(1-\mu)\bar{z} \in S$.

From our hypothesis C is a convex set; given any $x \in C_0$ and $\overline{y} = (\overline{y}^0, \overline{y}^1, \overline{y}^2)$ with $\overline{y}^0 \ge f_0(x)$, $\overline{y}^1 \ge g(x)$ and $\overline{y}^2 = h(x)$, we must find a $\mu > 1$ such that $\mu y + (1 - \mu) \overline{y} \in C_0$.

By convexity and $x^0 \text{eri}(C_0)$ there is a $\lambda>1$ such that $\mu x^0 + (1-\mu)x \in C_0$ for all $1<\mu \le \lambda$. Since $h(\cdot)$ is linear-affine

$$\mu y^{2} + (1 - \mu) \bar{y}^{2} = \mu h(x^{0}) + (1 - \mu) h(x) = h(\mu x^{0} + (1 - \mu)x)$$
(1)

and since both g(') and $f_o(\cdot)$ are continuous on $ri(C_o)$ the definition of y^o and y^1 shows that

$$\mu y^{0} + (1 - \mu) \tilde{y}^{0} > f_{0} (\mu x^{0} + (1 - \mu) x)$$
 (2)

and
$$\mu y^{1} + (1-\mu)\bar{y}^{1} > g(\mu x^{0} + (1-\mu)x)$$
 (3)

for some $1 < \mu < \lambda$. (1)-(3) imply that $\mu y + (1-\mu)\tilde{y} \in C$ so that $y \in ri(C)$.

Theorem 3: Theorem 1 and Theorem 2 are equivalent.

<u>Proof</u>; (Theorem 2 \Rightarrow Theorem 1) Letting $C_0 = C_1 \times C_2 \times \mathbb{R}$ and $x = (x^1, x^2, x^3)$ the Fenchel problem is restated as

$$v = \inf_{x \in C_0} \{f_1(x^1) - f_2(x^2) : x^1 = x^3, x^2 = x^3\}$$

This is an ordinary convex problem with the associations $f_o(x)=f_1(x^1)-f_2(x^2)$ and $h(x)=\binom{x^1-x^3}{x^2-x^3}$. Its Lagrangian dual problem is

$$\frac{1}{d} = \max_{\alpha j \in \mathbb{R}^{n}} \inf_{x \in C_{0}} \{f_{1}(x^{1}) - f_{2}(x^{2}) + \alpha^{1}x^{1} + \alpha^{2}x^{2} - (\alpha^{1} + \alpha^{2})x^{3}\}.$$
(4)

The hypothesis $\operatorname{ri}(C_1) \cap \operatorname{ri}(C_2) \neq \emptyset$ implies the Slater condition for this problem and consequently that $v=\overline{d}$ by Theorem 2. Since $x^3 \in \mathbb{R}^n$, the infimum in the dual problem is $-\infty$ whenever $\alpha^1 \neq -\alpha^2$. Therefore the dual maximization occurs for some $\pi = -\alpha^1 = \alpha^2$ and the Lagrangian dual (4) can be written as the Fenchel dual, giving $d=\overline{d}=v$ to prove Theorem 1.

(Theorem 1 \Longrightarrow Theorem 2). Letting C_1 be the set C of Lemma 1 and letting

 $c_2 = \{(y^0, y^1, y^2) \in \mathbb{R}^{1+m+r} : y^1 \le 0 \text{ and } y^2 = 0\}$ the Lagrange problem is written as

$$v = \inf \{y^o: y = (y^o, y^1, y^2) \in C_1 \cap C_2 \}$$
.

If x^o is the Slater point of Theorem 1 and $y^o > f_o(x^o)$, $o > y^1 > g(x^o)$, $y^2 = h(x^o)$, Lemma 1 implies that

$$(y^0,y^1,y^2)$$
eri (C_1) $).$

Identifying $f_1(y)=y^0$ and $f_2(y)=o$, the Fenchel dual to this formulation is

d=max {
$$\inf_{y \in C_1} (y_0 - \pi^0 y^0 - \pi^1 y^1 - \pi^2 y^2) + \inf_{y \in C_2} (\pi^0 y_0 + \pi^1 y^1 + \pi^2 y^2)$$
}

and by Theorem 2, v = d.

Since the second infimum is $-\infty$ if $\pi^0 \neq 0$ or $\pi^1_j > 0$ for some $j=1,\ldots,m$ and is o otherwise, this dual problem reduces to

$$\begin{array}{ll} \overline{d} = \max & \inf \left\{ y^{O} - \pi^{1} y^{1} - \pi^{2} y^{2} \right\} = \max & \inf \left\{ f_{O}(x) + \pi g(x) + \alpha h(x) \right\} = d \\ \pi^{1} \leq o & y \in C_{1} & \pi \geq o & x \in C_{O} \\ \pi^{2} & \alpha & \end{array}$$

so that v = d to prove Theorem 2.

Observe that the sets C_1 and C_2 used above are those typically constructed in a separating hyperplane approach to Lagrange duality. Also, the approach utilized here can be applied to relate other versions of the saddlepoint theories, for example when the dual problems do not have optimizing solutions. Other extensions can be obtained. For instance, the problem

$$v = \min \{ \sum_{j=1}^{k} (A^{j}x) : A^{j}x \in C_{j} \}$$

where C_j are convex sets, A^j are real matrices with n columns, and f_j is convex on C_j can be expressed in Lagrange form as

$$\min_{\mathbf{x} \in C_0} \begin{cases} k \\ \sum f_j(A^j x^j) : A^j x^j = A\bar{x} \end{cases}$$

where $C_0 = C_1 \times C_2 \times \ldots \times C_k \times \mathbb{R}^n$ and $x = (x^1, \ldots, x^k, \overline{x})$. Writing the Lagrange dual and simplifying provides the dual problem

$$d = \max_{\pi j} \begin{cases} k & k \\ \sum_{j=1}^{k} \inf_{x \in C_{j}} (f_{j}(x) + \pi^{j}x) : \sum_{j=1}^{k} \pi^{j}A^{j} = 0 \end{cases}$$

and v = d if there is an $x^0 \in \mathbb{R}^n$ with $A^j x^0 \in ri(C_j)$. When r=2 and A^l is the identity matrix this is the dual problem of Rockafellar [5]. These duality correspondences can be carried even further by appending inequality and equality constraints to the above formulation.

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