AN ARCHITECTURAL COMPARISON OF CONTEMPORARY APPROACHES AND PRODUCTS FOR INTEGRATING HETEROGENEOUS INFORMATION SYSTEMS

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ABSTRACT

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ABSTRACT

Virtually all large organizations are faced with the problem of intelligently accessing and integrating information maintained in islands of disparate computer-based systems. Since a number of vendors claim that their respective products and services can overcome the problem of integration, a comprehensive study was made into the technical characteristics of all these products and services, as well as additional ones identified through other means. This report documents the results of this study in terms of a new taxonomy for classifying various approaches for integrating heterogeneous information systems. In addition, the principal concepts and issues pertaining to this field are also presented in this report.

SECTION ONE: INTRODUCTION

1.1. Background

Most organizations today depend on a portfolio of information processing machines, ranging from mainframes to microcomputers and from general purpose workstations to very specialized systems, to meet their computational requirements. While the individual hardware and software components in the growing array of computer systems may meet the objectives for which each was initially designed, their heterogeneity presents a major obstacle to ready access and assimilation of information. It requires tremendous amounts of time and effort to retrieve information from multiple systems today, and in many situations, it is simply not feasible to obtain integrated responses in desired timeframes.

As compared to the intra-organizational situation discussed above, the issues involved in integrating information on interorganizational basis are even more complex. First, there is greater heterogeneity in technical capabilities across organizations as compared to a single organization. Second, there is less likelihood of reaching consensus on common goals and standards. Third, non-technical issues, such as access rights and sharing of costs, acquire added importance when multiple organizations are involved. Integration of relevant information from various computer-based systems is indeed a challenging task.

1.2. Functional Deficiencies

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The problem of inefficient, incomplete, and time-consuming access to information can be traced, from a technical viewpoint, to functional deficiencies at several levels as described below:

Structured and Unstructured Applications

Conventional and computer-based information systems have been designed with specific applications in view. These systems are efficient for performing the originally intended application, but they are inefficient in dealing with ad hoc queries or new unanticipated applications.

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Information versus Knowledge

Traditional database systems focus on retrieval of data, and on performing elementary operations (e.g., sorting and merging). Such systems cannot use both data and programs to respond intelligently to queries like "How many vehicles can be made available to evacuate personnel from the vicinity of a disaster site within a specified period of time?" which require analysis of many factors.

Diverse types of Information

Different types of information (numerical, graphic,

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pictorial, speech, and video) are referenced in very different ways. During the sixties, computers were designed to manipulate, store, and retrieve numerical information. During the seventies, the focus was on textual information. Relatively less effort has so far been directed towards efficient storage and retrieval of pictorial information and its combination with numerical and textual information.

Communications

In spite of continuing effort by a number of national and international institutions, it is still difficult to transfer complex information across computers of dissimilar architectures.

<u>Granularity</u>

It is difficult to judge what volume of information should be made available, and how it should be arranged and tailored to meet the needs of the person requesting the information.

Security

It is necessary to determine who can access the information, and what each person is permitted to modify.

This aspect is especially important for applications involving more than one organization.

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Semantics

It is difficult to specify exactly the subset of information desired without knowing details of the systems being accessed. Further, a particular piece of information may possess different connotations on various systems. Conversely, the same item (such as a particular spare part) may be specified by different numbers on various systems. The above list describes a few of the key technical problems.

In addition, there are major non-technical impediments that restrict integration. Because of the large number of issues and variables, a global solution to the problem of integrating distributed information systems has remained an elusive dream. Partial solutions are becoming available, however, and these solutions are identified and analyzed in later sections of this report.

1.3. Levels of Integration

Integration of information resources can be visualized at different levels. At the least complex end of the spectrum, information can be exchanged across a network using basic facilities such as remote

procedure call [13] and message passing. At the next higher level, a new layer of software can be created to insulate users from idiosyncrasies of different host machines. This layer can provide uniformity at the level of logically centralized file servers as in the case of Andrew at CMU [2], or at the level of application programs as in the case of Project Athena at MIT [3].

The approaches described in the preceding paragraph require significant degree of user involvement. Also, they represent instances of integration of information. The focus of this report is, however, on integration of information systems. Further, the emphasis is on systems "dissimilar" from each other in terms of:

- *Hardware, for example, one information system using an IBM mainframe and a second system functioning on Unisys hardware;
- *Operating systems, for example, one information system using OS/MVS, and other using UNIX;
- *Data models, for example, one system using network data models, and the other using relational data models;
- *Database management systems, for example, one using INGRES, and the other using ORACLE database management systems, respectively.

Apart from the above differences arising from dissimilarities in hardware and software, different information systems use different sets of design decisions and underlying assumptions. For example, one system may be using annual salary of personnel while another may be using monthly salary; further, the salary in one case may be gross salary (typical of budgeting systems), whereas in the second case, it may be net salary (typical of systems used for printing checks). Similarly, a particular spare part may be

specified by one part number in one system and by a different part number in a second system. All these differences must be reconciled when attempting to integrate dissimilar information systems.

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There are two facets to integration--physical integration and logical integration. <u>Physical integration</u>, accomplished through physical connectivity, refers to the process of establishing actual communication links among disparate systems. <u>Logical integration</u> refers to the process of accessing disparate systems in concert for generating integrated responses. This involves reconciliation of different assumptions and perspectives embedded in the systems being integrated. At the logical level, the integration system must:

- *Know where all the information is stored, along with the data formats and the query languages of the local systems;
- *Decompose the query into subqueries that can be executed by local systems;

*Accumulate the results from all the subqueries;

*Reconcile differences among the results accumulated; and

*Formulate integrated answers.

All the above tasks must ideally be carried out with no modification to existing systems. However, many products discussed later in this report require significant changes to existing information systems.

Based on the level of integration achieved, and some additional

characteristics, various types of integration approaches can be distinguished. These approaches, using the terminology of different researchers who have used dissimilar terms to refer to same or similar ideas, include the following:

> *Cooperation System for Heterogeneous Data Base Management System [4].

*Federated Database [5,9].

*Composite Information System [6].

*Knowledge Based Integrated Information System [7,11].

*Multidatabase [8,10].

*Superdatabase [12].

Instead of attempting to cover all the views expressed by the individuals who coined these terms and phrases, the endeavor in this report is to encapsulate the main concepts and the emerging technologies and products that enable integration of dissimilar information systems.

1.4. Constraints

The design of an integrated information system is constrained by the following facts:

- (a) The new system must integrate existing information systems, each of which was designed independently with little premonition, if any, that the system would form part of a larger system at a future date.
- (b) Each existing information system must be fully understood in order to comprehend all underlying assumptions and semantics. Rarely are these systems fully and accurately

documented.

(c) The design must allow flexibility to add or remove individual systems into the integrated configuration. As such, it requires detailed knowledge of both immediate and long-term constituents of the integrated system.

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(d) The design must involve the least amount of changes to existing systems. Transactions, irrespective of whether they are queries or updates, involve retrieval or manipulation of information maintained in individual systems, each of which imposes its own set of requirements and standards which must be reconciled by the new system.

The goal is to mitigate all the above-mentioned problems in an efficient and effective manner.

1.5. Models and Schemata

A Distributed Heterogeneous Database Management System attempts to meet the above goal by providing the facilities to access, to aggregate, and to update information maintained in multiple, distributed, heterogeneous systems. A typical distributed heterogeneous approach is characterized by the following:

- a) The use of a common data model, with appropriate mapping to data models of participating systems;
- (b) The support of a standard user language for retrieving data on a global basis in all cases, and also for storing data in some cases;

- (c) The ability to decompose a user request into an equivalent set of subqueries expressed in the format and manner required by participating systems;
- (d) The capability to resolve incompatibilities in semantics,data structures, and other parameters;
- (e) The mechanism to take the individual pieces of information provided by participating systems and to integrate them in order to provide a cohesive response to the original user request.

In addition to the above, some approaches facilitate implementation of sophisticated features such as authorization control and the ability to draw inferences in situations involving missing or conflicting data.

When data need to be retrieved from a single file, it is sufficient to simply know the file design including the keys for accessing records. One single view or schema of the file design meets the requirements of the systems analyst and the user. However, when there are several disparate information systems and there are many users from different disciplines, each user tends to look at the information from his or her perspective. By having more than one schema, it is feasible to handle complex problems, as well as to achieve data independence.

Most experts recognize the requirement for three types of schema. This notion of a three-schema architecture originated in 1977, when the Study Planning and Requirements Committee (SPARC) of the

American National Standards Institute (ANSI) published a report on this subject [14]. The three types of schema are as follows:

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- (a) <u>Global conceptual schema</u>: This schema is intended as a blueprint that shows the existence of all the information maintained in all the information systems. To enable intelligent access and integration, the conceptual schema must be designed to incorporate the semantics of all stored data. For this reason, it is sometimes referred to as logical schema. There is only one global schema for the entire enterprise.
- (b) Local or internal schema: This schema covers the data in an individual information system, that may itself be hosted on one or more computers. Since there is one internal schema for each information system, a complex environment necessarily implies a large number of such local data models. The collection of such local models is referred to as a local schemata.
- (c) External or application schema: This schema describes the user or the application view of the distributed data. Since each class of users or applications will view the data in a different manner, more than one external schema is required. This collection is referred to as external schemata.

The strategy used to map information between these three types of schemata determines the overall capability to mediate across dissimilar architectures, to maintain integrity of data, and to meet user expectations. Transformation of information between the

three schemata is ideally managed via a three-schema data dictionary. Use of very powerful data definition language is critical in defining the global schema in terms of objects, events, and integrity constraints [15].

It is important to emphasize that many commercial products and approaches do not support the concept of a global schema. Instead, a relational table or matrix is established between the required external schemata and the available internal schemata. Such a twoschema approach, referred to as the interfacing approach, offers no technological control over data integrity [15]. Data are exchanged either through transaction passing or by controlled redundancy. In the former case, all transactions are transmitted from the source database to the destination database, where they are transformed to adhere to the requirements of the external In controlled redundancy, instantaneous copies of the schemata. source database are stored at multiple nodes of the network, and these copies are replaced by new versions at periodic intervals. Both alternatives of the interfacing approach attempt to meet new user and application needs in an ad hoc manner, without any consistent set of rules governing the overall computational environment.

To overcome the above problems, it is necessary to use a threeschema approach, especially in applications of a critical or a strategic nature. All the major research programs funded by the U.S. Government, such as Multibase, Integrated Information Support

Integrated Manufacturing Distributed System, and Database Administration System, adopted the three-schema approach. Comparative evaluation of eight major initiatives in this area appears in [16,17]. To maintain adequate integrity of shared data of critical and strategic importance, the need for utilizing a three-schema approach cannot be overemphasized. This aspect acquires added importance when an application needs to access data from various DBMS-based information systems, as well as from earlier file systems that contain no notion of DBMS.

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1.6. Languages

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The only database language that has gained endorsement from major computer manufacturers, and from standards establishing organizations such as ANSI and ISO, is SQL. SQL statements can be used to create tabular views that approximate the notion of an external schema. Further, it can be used to create a conceptual schema, although a semantically weak one [18]. Intended to be used for queries, SQL has no internal schema.

There are several flavors of SQL in existence. A number of vendors claim to offer support for multiple versions of SQL. Further, some SQL based integration approaches can deal with legacy systems that offer no support for SQL. In such cases user queries expressed in SQL are converted into the native languages of the participating information systems.

There is a recent trend towards using graphics techniques to select databases, files, and tables. From a user perspective, these iconoriented approaches are far friendlier than traditional language oriented techniques. However, there is still no common graphics interface that has been endorsed by a critical mass of vendors. Even with a proprietary graphics interface, navigation with icons may be especially relevant for users at the top echelons of the organization.

1.7. General Direction

The basic concepts outlined above have been used in evaluating the principal products and approaches that are commercially sold. The list of vendors contacted is placed at Appendix A, and detailed specifications are tabulated in Appendix B. For the subset of products and approaches that hold some degree of relevance, a brief write-up was generated on each product. These write-ups are presented in Section 2 of this report. In Section 3, a new taxonomy is created for classifying different approaches to integration, and recommendations are made for selecting an appropriate set of technologies.

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SECTION TWO: HIGHLIGHTS OF PRINCIPAL PRODUCTS AND APPROACHES

In Section 1, the fundamental principles of integrating dissimilar information systems were presented. In this section, principal products and approaches are discussed and analyzed using these fundamental principles.

The products and approaches described in this section were derived from numerous sources, including articles in computer journals and magazines, proceedings of relevant conferences, lists of conference participants, and personal discussions with representatives in various companies. From the initial list generated, the systems previously covered in References [1] and [16] were deleted, unless the complexion of the system had changed significantly during the intervening period. The material presented in this section is not an exhaustive survey of all issues, but rather a discussion that covers the whole field taking all significant design approaches in view. It should be emphasized that some companies offer customized services, rather than a product, making it difficult to evaluate them in detail. In this section, each product is described by an introduction, its functional architecture, and its salient features. Additional details are tabulated in Appendix B.

2.1. BBN Systems and Technologies: Cronus

Cronus is an environment designed to support coherent integration of heterogeneous computer systems. The objectives of Cronus are: to provide an effective, general-purpose, distributed computing environment composed of hosts with differing hardware and operating systems; to support an integrated computing environment consisting of both the native operating system environment and the Cronus environment; and to provide comprehensive support for development of large-scale distributed applications. Cronus began as a research project funded by the Rome Air Development Center and the Naval Ocean Systems Center in 1981. The project was a two-year joint effort between BBN Systems and Technologies Corporation and Xerox Advanced Information Technology, both of Cambridge, Mass. The first prototype from this project became operational in 1984.

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Cronus supports both object-oriented and relational local data models. Global updates may be performed on an object by object basis. Update operators for insertion, deletion, and modification are in algebraic form and operate associatively on sets of tuples. Stored update procedures may include conditional statements and expressions.

FUNCTIONAL ARCHITECTURE

Cronus runs on top of existing operating systems on various hardware platforms. The Global Schema definition is supported for object-oriented data storage and retrieval. Every system resource

is an object and is accessed through operations defined by the type of the object. Cronus supports heterogeneity by serving as a bypassable layer of abstraction between applications programs and native operating systems. Through this approach, both conventional access to native operating systems resources and services and access to a coherent, uniform (object-oriented) system interface, regardless of computer system base, are supported.

Distributed query processing and update capabilities require translation of high-level requests into sequences of algebraic operations on data stored at possibly different sites. The algebraic approach defines a collection of "physical" operators on tuples and takes into account the sort order, the presence of indexes, and the physical location. Transformation rules determine when two algebraic expressions are equivalent, and optimization rules specify when it is profitable to transform an expression into another equivalent one.

Salient Features

Cronus incorporates an evolutionary approach in its design. Cronus claims that its Open Architecture approach makes it possible to upgrade or replace components or technology with little or no impact on the rest of the system. There is uniformity across diverse operating systems, hardware platforms, and applications.

2.2. Booz Allen: IDIMS

IDIMS is an expert system oriented approach designed by Booz Allen to provide automated data access. The user obtains access throughout the distributed data environment with a single global manipulation language and a user-friendly, intelligent interface. Upon receipt of a query, IDIMS decomposes the query into a set of database calls and function applications; determines security violations on databases; formulates queries in appropriate languages and queries databases; applies functions to retrieve data and finally displays the data via a state of the art graphic interface, in a style appropriate to the type and the number of data elements to be displayed.

It is important to note that IDIMS is not a commercial "off the shelf" product, but instead a conceptual framework used by Booz Allen to achieve automated data access. Although a complete production system has not yet been implemented, Booz Allen claims that it has built systems which relate to each functional component of IDIMS.

FUNCTIONAL ARCHITECTURE

The IDIMS system consists of a Query Processor, a Knowledge base for database models and function models, a Query Language Generator, and a Natural Language Processor.

The Query Processor translates a global data manipulation language

query into a series of data and function calls using the database and function models knowledge bases. In doing so, it plans the appropriate databases and functions to apply, chooses between plans using query optimization and time consideration techniques, and executes plans by calling the Query Language Generator to retrieve data. In addition, it determines whether there are potential security violations for the local databases queried.

The Data Base Models Knowledge Base is a structure which maintains data models for each of the distributed local databases. These models contain information and rules on security consideration, representative query language, data, data type, and host computer.

The Function Models Knowledge Base is a structure which maintains knowledge about the available procedures to alter or determine data attributes. Each function model provides for a specific logical data linking procedure such as correlate, sum, and join, and describes domain, range, and process definition. The Query Processor uses the Function Models Knowledge Base to determine which functions satisfy given constraints.

The Query Language Generator translates the global manipulation language into the appropriate local database query language. It then makes these queries and retrieves the data. Accordingly, the Query Language Generator maintains translations for each local database language and procedures to send queries to each computer on the network.

The Natural Language Processor provides an additional interface through which individuals can use limited natural language capabilities to query databases. This processor then translates the natural language into the IDIMS global manipulation language using conceptual dependency and case grammar.

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SALIENT FEATURES

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IDIMS is not a product, but an approach that eliminates the need for the user to know multiple query languages, locations of data, database paradigms, and manipulation schemata. A high-level, iconic, global query language is supported; however, updates are not supported.

2.3. Cambridge Technology Group: SURROUND

The Cambridge Technology Group (CTG) of Cambridge, Mass., specializes in providing customized distributed heterogeneous data integration within short time frames. A small group of users and managers from the client company interface with SURROUND developers at CTG to implement software interfaces as well as a rapid prototype of a functional application within a one week period. The goal of CTG is to construct a user-friendly system which focuses on one particular strategic application.

CTG applications support queries and updates. Local data models are supported via SURROUND interface and coordination tools. Complete data transparency is provided for geographically dispersed

locations via dialup ASCII modems.

FUNCTIONAL ARCHITECTURE

The UNIX operating system is used as a major part of the SURROUND service to resolve incompatibilities of different systems. Software tools access various distributed environments, extract information from individual programs, and then transfer data to the UNIX environment where the information is joined together. Talk Tools are communication tools which imitate the code that the host computer understands. The host computer accepts commands from the Talk Tool and sends results back to the same Talk Tool in the UNIX environment. There are also a variety of tools for attaining connectivity, security, network management, and for gauging performance.

The construction of the global schema is tailored to the needs of the application. SURROUND does not define a data manipulation language; it supports the interface currently used by the user. SQL formatted files are accessed with modified versions of SQL Data Handler and communication tools on remote system perform global updates. Commander, a proprietary tool, handles transaction management functions.

SALIENT FEATURES

Rapid prototype systems have been developed for TRANSCOM, General Electric, and other organizations within one week spans of time.

CTG claims that it can provide production software in as little as three months. Record structures can be changed easily and new databases can be added quickly, making the SURROUND environment flexible.

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2.4. Cincom: SUPRA

SUPRA, a product of Cincom, provides a commercial implementation of the ANSI/SPARC Three Schema Architecture. Such an architecture, described in Section 1.5, insulates applications from changes in the physical storage and conceptual database, reduces development and maintenance costs, and increases the useful life of applications.

FUNCTIONAL ARCHITECTURE

The SUPRA system integrates distributed heterogeneous databases and manages distributed processing by having a SUPRA kernel reside locally within each host machine or database. This has the effect of making the hosts aware of the integrated database network and in essence creates a "standard protocol" through which the databases can communicate.

The SUPRA kernel consists of two basic engine components, the Distributed Relational Data Manager (DRDM), and the Heterogeneous Data Management Processor (HDMP), along with a portion of the Global Directory. Using the kernel on each local host, the requesting node coordinates the entire unit of work, including the

decomposition of the query into separate neutral protocol requests and the determination of where to send them. The local/receiving nodes, each with its own kernel, translate the request into the appropriate local data manipulation language, gather the data, perform calculations, and return the result to the requesting node.

In this decentralized structure, each node is intelligent and has the ability to coordinate with each other via knowledge of the network in the resident Global Directory. This approach to distributed database and processing management assumes that most of the data accessed by any given node reside locally and that an "independent" node, with the ability to process locally, has significant speed and reliability advantages. In those instances where the data are not local, the Global Directory contains the information to tell the DRDM where to get them. As such, the network is not dependent on the reliability of a central Global Directory or engine. If any one node goes down, only the data residing at that node become inaccessible, not the whole system. Additionally, by distributing the necessary processing information, every transaction need not access a central directory, thus eliminating a potential bottleneck. However, with more than one copy of the Global Directory, the SUPRA approach implies more maintenance and overhead compared to a single copy approach. Every time the Global Directory is updated at one node, the update must be reflected across the entire network creating a "cascade" of transactions throughout the network.

SALIENT FEATURES

SUPRA currently provides for the use of synonyms and the resolution of scale conflicts. The decentralized approach allows a synonym table oriented to each user's particular application and working language. The Global Directory contains the ability to retain rules for performing scale translation and calculation before data integration. A unique feature called "triggers", alerts the user to the possibility of a data conflict. For example, a trigger could be established to ensure that two numbers, before they are combined together, be within a certain percentage of each other. At execution time, when a user wishes to combine the two numbers, a violation of this requirement would lead to an interrupt, and the user to be prompted about the problem.

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SUPRA utilizes a two phase commit procedure which provides good concurrency control. While assuring secure updates for SUPRA or other database management systems which support two phase commits, the product currently only allows updates from join views to distributed databases which use SQL.

2.5. Cognos Corporation: POWERHOUSE

PowerHouse consists of four components: PowerHouse 4GL, PowerHouse StarBase, StarNet and StarGate. PowerHouse 4GL is a fourthgeneration language; PowerHouse StarBase is a relational database management system; StarNet is network routing software; and StarGate is a gateway provider to other supported local

database management systems. The PowerHouse 4GL was developed by Cognos Corporation, Ottowa, Ontario while the remaining components were developed by InterBase Software Corporation, Bedford, MA. Cognos began to market the Relational DBMS of InterBase and its distributed technology in conjunction with the PowerHouse application development language in 1988.

FUNCTIONAL ARCHITECTURE

In this approach, there is no concept of a single global schema. Instead, multiple relational schemata may be referenced, even from within the same program. Relational operators are utilized to selectively access and merge pieces of data. Research on a distributed query engine is currently in progress.

A discussion of the capabilities of InterBase components is provided later under its own prototype overview. The main difference between PowerHouse and InterBase is global data manipulation language supported. Both products support SQL and Groton Data Manipulation Language, but the stand-alone PowerHouse 4GL is unique to PowerHouse.

One of the benefits of the PowerHouse 4GL is that it provides access to relative Record Management Service files and network databases in addition to sequential and indexed Record Management Service files and relational databases supported by InterBase. The Open Systems Relational Interface (OSRI) is a relational database protocol. PowerHouse 4GL elements for transaction

processing, report generation and batch updating use a relational access approach which makes it possible to embed OSRI calls into the elements. This means that as more local data models are added to the configuration, it is possible to develop applications which access them using PowerHouse.

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The PowerHouse 4GL product set includes PhD master data dictionary. Using PhD, it is possible to treat all meta-data as if they were resident in a single logical dictionary. Thus, one need not redefine meta-data in the PhD dictionary.

SALIENT FEATURES

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The salient features listed in the description of the InterBase prototype also apply here. Cognos is attempting to incorporate expert systems and object-oriented techniques in the PowerHouse product.

2.6. Control Data Corporation: ASCENT

ASCENT is derived from the Technology Information System (TIS) developed by Lawrence Livermore National Labs (LLNL). Work on this system originally commenced in the early seventies. CDC has a technology transfer license with LLNL (actually with the University of California) to support this technology, and to take it into production.

ASCENT is a suite of software products that provides

communications, networking, office automation, and database services within an interactive environment that is common to all users. The *gateway portion of ASCENT supports a powerful interpretive scripting language that can simulate an interactive user when connected to a heterogeneous remote computer. This gateway serves as a tool that can be programmed with the scripting language to log into databases and to extract data from a broad range of remote systems.

FUNCTIONAL ARCHITECTURE

A collection of software modules, known as the Intelligent Gateway Processor (IGP), use a UNIX system as the link between menu type front-end systems and geographically and technically disparate mainframes. The IGP front-end program, Integrated Information System (IIS), offers menus composed of descriptor files which might consist of classes of information resources, specific connections to remote hosts, local database options, or available system utilities. In addition to menus, IIS understands a command set that overlaps with the UNIX Shell command set. As IIS is implemented in a script-driven interpreter, UNIX commands can be recognized even without arguments, and prompts can be enabled for arguments or options.

The Network Access Machine (NAM), an extended implementation of the specifications published by the National Institute of Standards and Technology, utilizes a variety of communication methods to provide access to other computer systems. Currently NAM can

interface with direct asynchronous connections between two hosts, most autodial modems, TCP/IP networks and X.24 networks. NAM uses an interpreted communication language to specify the logic necessary to establish connections and conduct dialogues with remote systems. This language includes structured constructs, extensive pattern matching capabilities and capabilities to insert "filters" in an existing connection. These filters are specialized programs written to modify the appearance of the data.

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In a heterogeneous environment, ASCENT*gateway must be driven by the "host UNIX system" database management system to extract specific data elements from their remote computers under script language control. The DHDBMS places these data elements into the proper place in the "logical record" requested by the user. This approach requires no software to be added to any of the remote systems since each system feels that the scripting language interactions with other remote systems are identical in appearance to interactions with its own users.

SALIENT FEATURES

ASCENT*gateway allows automated access and connectivity to local or remote computing resources. ASCENT*mail provides complete electronic mail and message services within a computer network including addressing, forwarding, and mail file management functions. A menu-driven user interface is available to access integrated software applications, utilities, and system administration tools.

2.7. Control Data Corporation: CDM*PLUS

The concept of CDM*plus was developed and prototyped with direction and funding from the Integrated Information Support System (IISS, pronounced I^2S^2) initiative of the U.S. Air Force. CDM*plus defines and categorizes data using the ANSI/SPARC Three Schema Architecture discussed in Section 1.5.

FUNCTIONAL ARCHITECTURE

The CDM*plus dictionary/directory is loaded and maintained via the Neutral Data Definition Language (NDDL) which describes all metadata to the ANSI/SPARC based three schema structure. NDDL is SQL based and describes objects with verbs such as create, define, alter, and drop. Each schema has unique commands. For example, in the internal schema, the NDDL can describe the three basic types of database models. NDDL provides inter-schema mappings using the conceptual schema or the data model as the base. NDDL is a human and machine readable language that can be used in either in interactive mode or in batch mode, as well as on forms.

The mapping strategy between the schema allows data to be accessed without regard to the location or the storage mechanism. For example, an application can join two relational user views from the external schema that each map to two entities. These entities in turn can be mapped to multiple databases and database types on multiple computers without burdening the user or developer with all of the details of the actual conversions, translations or mappings.

This independence is critical to provide a buffer for life cycle maintenance as applications and storage structures change, and to provide the capability for full relational operation (including outer joins and referential integrity) between heterogeneous databases.

III

Queries can be embedded into current applications written in Cobol, C, or Fortran. These queries are precompiled by CDM*plus using the external schema view and transformed into actual database queries. All code necessary to access the databases is generated including DBMS precompilations. Since the base schema is the conceptual schema or data model, domain constraint checking and referential integrity checking can be automatically enforced based on the conceptual schema. After the code is generated, it is transferred by the Code Manager to the appropriate computers where local compiles, DBMS precompiles, and links are accomplished.

All query processing is controlled by two CDM*plus components: the Network Transaction Manager (NTM), which controls unique process to process communication capabilities that are necessary in the distributed environment, and the CDM*plus Distributed Request Supervisor (DRS), which controls the actual queries on each host and formulates the messages sent to the NTM. This approach allows serial or parallel processing of queries. In the serial process, one query produces input for other queries.

Each local node contains an NTM and multiple CDM*plus runtime (CDMR) modules to facilitate distributed processing. The CDMR

incorporates DRS, Aggregators, File Utilities and Code Manager. Query modules in applications that attempt to integrate information from multiple heterogeneous databases can be replaced by NDML statements. This allows for easier maintenance of legacy applications because changes to the systems only require reprecompilation to generate all the necessary code.

SALIENT FEATURES

CDM*plus requires the user to model his or her enterprise and to load that model into the CDM. The three schema architecture provides one global conceptual view of the complete data. In addition, it provides independence between the actual physical storage of the data, the users' view of the data, and the conceptual view of the data. Modeling and integration of applications and databases with CDM*plus is done in an evolutionary manner, allowing for the implementation of one application at a time rather than a full rewrite of the entire system.

2.8. Data Integration, Inc.: Mermaid

Mermaid, a trademark of Unisys, is an integrated data access product marketed by Data Integration, Inc. It allows users of multiple relational or network database management systems, running on different machines, to access and manipulate data using standard SQL, even when the participating systems don not support SQL. Integration may require translation of the data into a standard data type, translation of the units, combination or division of

fields, and encoding values. Only updates to a single database are presently supported. Updates to replicated and fragmented relations may cross databases, but the updates to all databases are not made concurrently.

The Mermaid system is designed for an environment in which there are many databases that are independently administered and where it is not feasible to standardize on one DBMS. With Mermaid, new applications may be developed either at the global level when distributed access is required, or above an existing old database. Local administrators grant access permissions to users of Mermaid and then Mermaid adds an additional level of access control.

FUNCTIONAL ARCHITECTURE

The major processes are: User Interface Process, Distributor Process, DBMS Driver Process, and Data Dictionary/Directory (DD/D). The User Interface Process contains the SQL parser, the query validator, the query library, and the report generator. The Distributor Process contains the optimizer and the controller. The optimizer plans query execution and the controller starts processes and coordinates execution. There is one DBMS Driver Process for each database to be accessed. It contains the schema translation, the language translation, and the DBMS interface. The Data Dictionary/Directory (DD/D) is a database that contains information about the schemata, users, host computers, and network. It supports three layers of schema definition.

A user input is either a command to execute within the user interface or an SQL query. The SQL query is parsed, validated, translated into Distributed Intermediate Language (DIL), and sent to the distributor process which performs the Distributed Execution Monitor function. The optimizer plans the execution by consulting the DD/D for the location of the relations. It sends commands and subqueries to the controller for transmission to the DBMS drivers. The driver reads a message which may be a command to initialize, terminate, execute schema and data of the query from the global form into the local form and then translates it from DIL into the local DBMS language. The first time that a relation is retrieved, the data are translated into the global form.

SALIENT FEATURES

Mermaid treats files as large objects. Qualification is done on structured data. The user selects qualified objects to view. The system then starts a retrieval process and a display process that are customized for the file type.

Communication between processes is done using standard protocols. The main processes use TCP/IP on an Ethernet. If TCP/IP is not available on the computer with the DBMS, another protocol may be used to communicate between the controller and the driver or between the driver and the DBMS.

Mermaid is written as a set of processes with well defined interfaces so that it can be customized for different environments.

It can also be embedded in or called by other products and it can use other products to process the output.

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2.9. Digital Equipment Corp.: Rdb/VMS,DBMS

Digital Equipment Corp. (DEC) offers two alternative methods for distributed data management in the VAX/VMS environment. One method is geared towards a decentralized environment with users having their own processing resources and connected through a Local Area The other method is directed towards centralized Network. distributed processing with a VAX cluster comprising two or more VAX processors, and the database located on a common disk. Each of these methods can be used with either of the two DEC data management systems: Rdb/VMS is a relational DBMS while Digital's Database Management System (DBMS) uses a hierarchical/network model. Further, VAXlink Client is a layered VMS software product that provides unidirectional IBM mainframe database access and extract capabilities; it works in conjunction with VAXlink Answer/DB Extractor to access IMS and VSAM data structures and to transfer extracted data to Rdb/VMS databases.

FUNCTIONAL ARCHITECTURE

VAX SQL is Digital's version of ANSI SQL and is used as an interactive and software development interface for relational databases that conform to Digital Standard Relational Interface (DSRI). Used to define, access, and update relational databases, it is characterized by an interactive data manipulation language,

a data definition language utility, preprocessors and dynamic SQL for COBOL, FORTRAN, C, and PL/I. DEC's Database Management System uses the Network Manipulation Language (NML), a standard for network models.

A Database Control System (DBCS) provides the connections between the application program and the database locations for the Data Base Management System. Application programs transmit information to the DBCS. The DBCS transfers all information requested by a call for data to the User Work Area (UWA). Information in the UWA is then used to update the database.

The VAX Distributor distributes data from any Rdb/VMS, VIDA (an access tool for IBM databases), or any other database conforming to DSRI. The commands of the VAX Data Distributor allow the user to specify the data to be distributed, the type of copy to be made, the target node, and the transfer and update schedules. Three kinds of copies are permitted by this tool: extractions, extraction rollups, and replications. These copies effect only the target database.

Digital's distributed, active, open-architecture dictionary is the VAX Common Data Dictionary/Plus (CDD/Plus). VAX CDD/Plus provides one logical dictionary system that is physically dispersed, thus eliminating the need to develop separate dictionaries to support proprietary tools or applications.

SALIENT FEATURES

Digital claims to provide one of the best relational database management systems for the VAX/VMS environment. VMS is a widely used operating system for development and execution of a wide range of applications. Apart from the above products marketed directly by DEC, a number of other vendors have opted to select VAX hardware as the host system for reconciling the idiosyncracies of individual components in a distributed, heterogeneous environment.

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2.10. Gupta Technologies, Inc.: SQL SYSTEM

Gupta Technologies, Inc. of Menlo Park, Calif., supports relational database networking. The SQL SYSTEM allows applications running on a PC to use data residing on minis, micros, and mainframes. The primary elements of the SQL SYSTEM are: SQLBase, SQLWindows, SQLTalk and SQLNetwork.

SQLBase is a distributed relational DBMS for PC networks. It can be used as a standalone DBMS, as a database server for personal computer based workstations on a network, and as a gateway to mainframe database management systems. The SQL language syntax of SQLBase is fully compatible with those of IBM's DB2 and SQL/DS. Gupta Technologies claims to be the first company to offer advanced concurrency, automatic consistency, and dynamic recovery for a distributed relational DBMS on a personal computer.

SQLWindows and SQLTalk allow the programmer or user to utilize SQL

without having to develop and compile programs. SQLWindows is an advanced applications generator which supports graphics, menus, dialogue and help boxes and mouse pointing devices. It contains a complete procedural language and provides exits to other programs. SQLTalk is an interactive data management program for accessing SQLBase via the SQL language.

FUNCTIONAL ARCHITECTURE

The SQLNetwork components enable mini, micro and mainframe SQL database management systems to appear like distributed database servers on the same network. The three components comprising the SQLNetwork are: the SQLRouter, the SQLGateway, and the SQLHost.

Requests for data are handled differently based on the location of the databases. The SQLRouter is the software that directs SQL requests to databases within the PC network. SQLRouter must reside on each PC and requires NetBIOS on each workstation. The SQLGateway routes SQL requests to databases outside of a single PC network. It runs as a memory resident program on the PC/AT and accesses a file containing the databases to which it can establish a connection at load time. This file also contains the physical node location of these databases.

SQLHost software resides on the host computer and transforms the host DBMS into a database server that accepts SQL requests routed to it by SQLGateway. It translates SQL commands into a format acceptable to the host DBMS and forwards the information for

processing. The results are then sent back to the network. A different version of SQLHost is necessary for each different SQL DBMS because of variations in SQL and application program interfaces.

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SALIENT FEATURES

SQLNetwork was designed to allow applications to take advantage of the capabilities of a personal computer, to transparently access databases that reside on diverse servers, and to serve many users with a large amount of data. The user is provided with full cooperative processing capabilities in DOS, OS/2 and UNIX environments. SQLNetwork provides access to DB2 and development is underway for SQLHost Software to support SQL/DS, Oracle, Ingress, and Sybase.

2.11. Honeywell Federal Systems, Inc.: Distributed Database Testbed System

DDTS is a testbed emphasizing modularity and flexibility for distributed heterogeneous database management. The project began at Computer Sciences Center (CSC) of Honeywell with the aim of designing and building a system that could be utilized for evaluating alternate approaches and algorithms for managing data stored in several databases on different platforms. Although DDTS is not a product, Honeywell utilizes this technology in specific contracts.

FUNCTIONAL ARCHITECTURE

DDTS supports a five-schema architecture. The internal schema is divided into the global representation schema, the local representation schema and the local internal schema. In addition, there is an external schema and a conceptual schema which follow the terminology of Section 1.5. A bottom-up approach is used to design a global schema for multiple databases. First, a local representational schema is built from the description of the local database. Next, all local representational schemata are integrated to form the global representation schema. Finally, the conceptual schema describing semantic integrity constraints to be enforced by DDTS is constructed.

The DDTS system architecture is based on a model of distributed computing in which a set of abstract processors execute sets of concurrent processes. Local and remote processes communicate via message exchange. DDTS consists of a set of Application Processors (APs) which control user interfaces and manage transactions, and Data Processors (DPs) which manage data.

The GORDAS Command Interface (GCI) enables multiple retrieval and update requests expressed in GORDAS, a high-level query language. Before a GORDAS transaction is executed, the GORDAS requests are translated to an internal representation for processing by the Translation and Integrity Control (TIC) module. The internal representation possesses features of relational algebra and relational calculus, and is similar to an internal representation

of SQL.

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Materialization and access planning (MAP) result in a strategy for processing distributed queries, updates and transactions. The response time (the time elapsed between the start of the first transmission and the time at which the result arrives at the required computer) and the total time (the sum of the time of all transmissions) are minimized with MAP.

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The Distributed Execution Monitor (DEM) performs distributed execution of a transaction as directed by MAP. Based on status information returned by the Local Execution Monitors (LEMs) during execution, the DEM decides to commit or abort a transaction. DDTS uses a two-phase commit protocol during transaction execution. To resolve concurrency control conflicts, the transaction is rolled back and then restarted.

The Local Operations Module (LOM) of each DP acts as an interface between the LEM and the local DBMS. LOM provides local translation and optimization of subtransactions that have been assigned to the DP by MAP. The LOM translates subtransaction requests in the internal representation to data manipulation requests in the language of the local DBMS.

SALIENT FEATURES

DDTS is a useful tool for evaluating approaches to the various problems associated with distributed, heterogeneous database

management.

2.12. IBM: DB2, SQL/DS, SAA, and others

IBM faces a challenge in developing its integrating approach. It commands a major presence in the mainframe computer market and in the personal computer market. It can either come up with design approaches that enable integration of information hosted on its own systems, or it can develop approaches that encompass non-IBM architectures as well. So far, the emphasis has been on the former alternative.

The IBM System Application Architecture (SAA) serves as a framework for developing consistent applications across three IBM computing environments: System/370, AS/400 and PS/2. SAA is a collection of software interfaces, conventions, and protocols, and includes support for distributed files and distributed databases that are resident within one of these environments, and conform to certain design guidelines.

IBM is developing a new relational Unix oriented approach that would integrate their AIX line (System/370, RT, and PS/2) as well as communicate with their SAA database line (System/370, AS/400, and PS/2)). Details of this approach have not yet been released by IBM.

At a lower level, IBM provides support for linking information from distributed relational (SQL/DS) and hierarchical (IMS) local

database management systems. In addition, its Data Extract product (DXT) can extract data from DB2 tables, SQL/DS tables, DL/I data sets, VSAM data sets, and physical sequential data sets, and convert them for loading into other DB2 tables, physical sequential data sets, and CMS files. DXT users can also write generic data interfaces for accessing non-IBM (OEM) databases on MVS.

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In October 1988, IBM upgraded DB2 and SQL for distributed environments. The new DB2 and SQL (called SQL/DS) use relational concepts to integrate information. In either case, because a central DBMS is not assumed, each location has all the capabilities offered by a local relational DBMS and may participate in an overall distributed network as well. Local functions are not affected by the role of the system in distributed processing. Flexibility is provided for deciding how data will be distributed and processed. The SQL data manipulation language can access remote tables or a user may work with a local copy. Also, data may be stored in one large table or split into several parts to accommodate data which may be accessed and updated only at certain local locations. Extracts, snapshots, and data replication are various types of copies which enhance retrievals and updates.

A user may extract a copy of one database table and load it into another database. Extracts are useful for migrating data from one data format to another and for maintaining read-only copies of a subset of data that rarely changes. Snapshots are read-only copies of tables that are automatically made by the system on a periodic basis specified by the user. These are useful for locations that

need an automated process for receiving updated information on an occasional basis. With the data replication support, any location may update any copy of a table, and these updates are sent to other locations automatically. This is effective for situations that require high reliability and quick data retrieval with few updates.

FUNCTIONAL ARCHITECTURE

Distributed relational data are accessed either on a remote unit of work or on a distributed unit of work. Remote unit of work permits access to a single, remote relational DBMS within a unit of work. A unit of work is a set of one or more related requests which must all complete successfully or not at all. An application program can query and update data from any number of tables residing in a single DBMS within a unit of work. In addition, application programs can access data in multiple database management systems one at a time, in separate units of work. SQL/DS and DB2 have remote unit-of-work capability as well as support for single SQL requests.

Distributed unit of work provides access to multiple relational database management systems within a single unit of work. An application program can read and update data from any number of database management systems within a single unit of work. However, a single SQL statement must refer to data from a single DBMS. Key features of distributed unit of work are: complete location transparency, coordinated system security mechanisms to determine whether the application is authorized to access the remote systems

and data, and the two-phase commit protocol to ensure termination consistency. DB2 supports distributed unit-of-work.

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SALIENT FEATURES

IBM states that its leadership in stand-alone database management products resulted from its advanced data security, data integrity, and data recovery capabilities. Its objective is to apply this expertise to the implementation of its distributed relational products. DB2 and SQL/DS have referential integrity which ensures the consistency of data values between related columns of two different tables.

2.13. Information Builders, Inc.: FOCNET

FOCNET, derived from FOCUS Network Products, is marketed by Information Builders, Inc., headquartered in New York City. Its nonprocedural language and DBMS is called FOCUS. Various FOCUS interfaces provide logical access to data in hierarchical, relational, and network models on remote machines.

FOCNET permits multiple individuals to update a database simultaneously. A change/verify protocol confirms that the record has not been updated by another user before it modifies the record. In cases where the record has been changed, the update is reapplied against the updated version. Also, this approach allows a user to define commit and rollback specifications for transactions.

FUNCTIONAL ARCHITECTURE

A client/server model layer is inserted on top of the operating system and the access-method architecture of each participating node. Its three components are client software, interfaces to network transport software and server software. This approach reduces network traffic and takes full advantage of available CPU resources.

Most local databases must have a Master File Description (MFD). All files must be described in the FOCUS Master Dictionary before FOCUS can access their data. The master dictionary provides much more than conventional file and data information. FOCUS file descriptions can also contain security restrictions, field aliases, data validation criteria, predefined relational join operations, and multiline column titles for use in reports.

FOCUS supports 4GL FOCUS, third generation languages, and ANSI SQL global data manipulation languages. Read/write interfaces generate the appropriate dialect of SQL for many well-known relational database engines which use variations from this standard data manipulation language.

Dynamic relational join operations permit linkage to up to 16 different tables and databases at one time. Processing is optimized by including record-by-record operations and set processing based on the functionality of the database. In addition, balance tree indexing accelerates the speed of data

retrieval and online updates of large databases. Users can logically invert FOCUS databases to further optimize retrievals.

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SALIENT FEATURES

FOCNET runs on many platforms under a variety of operating systems. The list of local database management systems supported is impressive as is the list of supported network services. FOCUS is the most widely used 4GL in the world. It is suited for developing software and provides for automated application development and maintenance.

2.14. Interbase Software Corporation: INTERBASE

Interbase is a product of Interbase Software Corporation, Bedford, Mass. Although Interbase operates in a variety of software and hardware environments, its Digital Standard Relational Interface (DSRI) compatibility makes it most suited for use with other DEC DSRI-compatible tools. Relational local data models, indexed and sequential Record Management Service files, and UNIX stream files are supported by Interbase.

Global updates to databases can be performed in the same interactive session, as Interbase allows multiple simultaneous transactions per process and provides for two-phase commits. These two features can be used to ensure that a transaction which updates more than one database will complete correctly or not at all. Partial commits and partial rollbacks are possible as well.

Interbase uses multigenerational data structures for data storage. Records are never updated, but only added. Old versions of the record are deleted by the next transaction which accesses the record. This means that update access to data being processed by another user is allowed. On actual update, Interbase will impose the record it originally retrieved with the record it is trying to update. If they are different, the transaction is rolled back and the user is informed of this fact.

Changes required to local databases are minimal. The Rdb database of DEC permits access to Interbase files through VMS. The ORACLE gateway requires an additional table to be defined on the database, access to tables in the data dictionary, and two servers to be started on the target node.

FUNCTIONAL ARCHITECTURE

Interbase supports two global data manipulation languages: SQL and the Groton Data Manipulation Language. These two languages may be used in precompiled third-generation language source programs or in the Query Language Interpreter (QLI). The QLI translates either language into BLR, a lower level language based on the DSRI, which may be generated and read by programs.

Although InterBase has no global conceptual schema, it provides a navigation path to distributed local database management systems through a database handler which defines node names for multiple databases. A program called Y-valve determines if requested data

are on a local node or a remote node. In the case of remote data, the request is routed through the network and the remote data dictionary is accessed.

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Interbase or RDB data requests are left in the BLR lower level language. With ORACLE, BLR is restricted to SQL through the gateway and passed to the ORACLE access method. Data format incompatibility between heterogeneous systems is transparently handled by a XDR layer.

SALIENT FEATURES

Interbase is designed to handle on-line complex applications (OLCP). It supports: multi-user, concurrent update; automatic transactions; serializable and high concurrency levels; two-phase commit; multiple database transactions; and multiple parallel transactions per process. These features provide excellent concurrency control capability.

2.15. Massachusetts Institute of Technology: Composite Information Systems

The concept of Composite Information Systems (CIS) has evolved at the Sloan School of Management of Massachusetts Institute of Technology. The first paper on this technique was published in 1978, and significant enhancements to the original concept have been made during the intervening years. While all other approaches discussed in this report tend to focus exclusively on technical issues, the CIS approach emphasizes that it is necessary to address

a number of inter-related strategic, organizational and technical issues, in order to implement systems that can effectively span applications, functional areas, organizational boundaries, and geographic separations.

Strategic issues include motivating cooperation between multiple organizations, each with its own goals, priorities, and security needs. One critical factor for such cooperation is participant consensus on the issue of access to each others' technical and nontechnical information. Issues such as domains of shared information, benefits to each participating group, and the role and the responsibility of each constituent must be resolved in advance, for the final system to succeed in a real world environment.

Organizational issues include the process of making controlled changes in complex organizational environments. The latest results from the realm of inter-organizational networks have been utilized to explain the multiple forces that modulate behavior of individuals, groups, and organizations. Also, the notion of focused standards has been conceived to serve as the foundation for neutral representation as well as for the delineation of more elaborated standards.

Technical issues relate to connectivity at both physical and logical levels. Physical connectivity refers to the process of actual communication among disparate systems, and covers aspects as bandwidths, security, availability, reliability, and inter-

network protocol conversions. Logical connectivity refers to the process of accessing disparate systems in concert for composite answers. One of the major challenges in logical connectivity is to reconcile the different assumptions and perspectives in the systems being integrated.

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FUNCTIONAL ARCHITECTURE

The operational prototype of CIS is called CIS/TK, which is abbreviated from CIS Tool Kit. This prototype is based on the three-schema approach, with one integrated global schema, multiple local schemata, and multiple application schemata. The CIS/TK architecture includes multiple Application Query Processors (AQPS), one Global Query Processor (GQP), and multiple Local Query Processors (LQPs). There is one AQP for each application, and this AQP converts an application model query, defined by an application developer, into a sequence of global schema queries, passes them to the GQP, and receives the results. Similarly, a Local Query Processor (LQP) interfaces between the global data processor and an individual query command processor (e.g., a DBMS).

The primary query processor is the GQP. It converts a global schema query into abstract local queries, transmits them to the appropriate LQPs, and joins the results before passing them back to the AQP. The GQP must know how to get the data, how to map global schema attribute names to the column names, and how to join results from different tables. This is done with the aid of concept agents. Conceptually, each concept agent is composed of

three functional components, as follows:

- (a) A concept definition which defines the rules and data, given the global schema and an application model;
- (b) A concept processor which pursues the goal of the concept agent, given its rules and data; and
- (c) A message handler which sends and receives messages on behalf of the concept agent.

The notion of concept agents is used to perform concept inferencing, in situations involving missing and conflicting information.

CIS/TK utilizes an object-oriented approach. Physical communication details and database idiosyncracies are encapsulated within the object. The object-oriented language utilized by CIS/TK is based on an enhanced version of Knowledge-Object Representation Language. CIS/TK is currently providing integrated access to several databases at MIT, and others outside MIT. External databases include Reuter's Textline and Dataline and I.P. Sharp's Disclosure.

SALIENT FEATURES

CIS is an integrated methodology that attempts to systematically deal with strategic problems, organizational problems, and technical problems that are inherent in any endeavor aimed at integrating islands of disparate information systems. CIS/TK is an operational prototype that utilizes innovative concepts to perform logical connectivity in an effective manner.

2.16. Metaphor Computer Systems, Inc.: DIS

DIS is a product of Metaphor Computer Systems, Inc. of Mountain View, California. Through the combination of a graphical interface and relational database technology, DIS lets users access data from multiple databases, perform analysis, and construct their own applications without programming. Designed for business professionals, DIS provides tools for strategic use of data from multiple sources; database query tools and report writers for direct and easy access to data; spreadsheets to analyze data; plot tools to view data graphically; and text processing tools to add commentary to the analysis.

Currently, DIS only provides for data access to SQL relational databases and to some flat files after they are converted to relational format used by DIS. There is no plan to provide direct access to non-SQL environments. However, to the extent local SQL database management systems can translate other databases into their own relational data model (e.g., Ingres), DIS can access them.

There is no logical schema which encompasses all database management systems in the network. However, a local schema of one or several databases is graphically depicted for the user. This provides the ability to "click" on each database and instantly see a list of its tables. DIS can be used to query multiple physical database management systems concurrently through multiple windows, and to combine and manipulate such data using the spreadsheet tool.

FUNCTIONAL ARCHITECTURE

DIS is a network of workstations (soon to include PS/2's and PC's) and servers that interact with relational databases. Databases can reside on Metaphor database servers as well as on non-Metaphor host machines. Metaphor files and communication servers can provide electronic mail and print services, plus access to a wide range of multi-vendor environments.

The main data access tool is the Query tool. When it is selected, the appropriate database is connected and each of its tables and its columns are displayed in a window. As the user chooses columns with the mouse, the Query tool draws a picture showing how the database tables will be joined to retrieve the data. Then it translates the graphic query into SQL and sends it to the database.

If data are not directly accessible to DIS via a gateway, they can be downloaded to a Metaphor database server. Host resident "formatters" are available for converting flat data files into the relational format used by DIS.

SALIENT FEATURES

This product claims to provide a true user friendly data access. Querying can be done without ever touching the keyboard. All tables and columns within a physical database are graphically provided to the user. The spreadsheet, the plotting facility, and the word processing tools are very practical features. The list of

government and commercial customers attests to the ability of this product to provide management with easy access to information. IBM has invested in Metaphor Computer Systems, Inc. and is currently working with it to provide a similar interface for IBM mainframes.

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2.17. On-Line Software International, Inc.: RAMIS Information System

RAMIS Information System from On-Line Software International, Inc., headquartered in Fort Lee, New Jersey, is used primarily for developing and running end-user applications. The system consists of many diverse components such as: a centralized database, a data dictionary, reporting and graphics facilities, links to external database management systems, data manipulation tools, a procedural language for application development, menu architecture development facilities, micro-to-mainframe communications, and data transfer facilities.

The RAMIS database may be comprised of both RAMIS and non-RAMIS files or databases. Most non-RAMIS files and databases may be accessed in a read-only mode. However, DB2 and SQL/DS relational databases may be queried and updated from within the RAMIS environment.

FUNCTIONAL ARCHITECTURE

RELATE is a facility to combine data from multiple files by operating on them as if they were flat relational tables. RELATE

provides relational capabilities regardless of the organization of the source of the data. The user must specify the database or the file, the data to be retrieved, the record selection criteria and how the data are to be associated.

The RAMIS Database Manager reads and updates RAMIS, DB2, and SQL/DS files. It allows multiple physical files to be logically interconnected to from a single data view. A user is presented with one relational view of the data. In the same way, multiple physical databases may be treated as one logical database. Specialized interfaces are required for accessing other non-RAMIS files and databases.

RAMIS has a totally integrated, menu-based architecture front end to all components of the system. Application developers can utilize the menu interface, use a command language, and construct their own customized menu interfaces.

SALIENT FEATURES

The RAMIS Information System operates on various IBM platforms under diverse operating systems. It provides a flexible environment for managing, retrieving, manipulating and reporting data stored both in RAMIS format and in external databases and files. Direct interfaces within RAMIS allow users to read non-RAMIS files and databases using RAMIS language facilities. It is not necessary to convert external files into a RAMIS file to access the data.

2.18. Oracle Corporation: ORACLE

ORACLE was developed by Oracle Corporation, Belmont, California, to serve as a general purpose tool for providing a uniform, integrated interface for retrieving data from several existing database management systems. ORACLE connects with SQL/DS, DB2 and DBase III relational databases. There is also an ORACLE utility for 123.

III

ORACLE is constantly evolving to meet System R* specifications which IBM determined in the 1970s to define the "ideal" database management system. The influence of IBM is further seen in the global data manipulation language which is a superset of SQL, the IBM developed Standard Query Language. ORACLE is a widelyinstalled DBMS in the VMS environment.

Both distributed query and single-node transactions are supported by ORACLE. For updates at multiple nodes, there is no assurance that transactions will either be committed or rolled back together. Data loss can only be prevented by application programming techniques. Global updates and the two-phase commit procedure will be supported in the next product release.

FUNCTIONAL ARCHITECTURE

The Open Systems Architecture of ORACLE for distributed processing and distributed database management is SQL*Star. It is comprised of three ORACLE products: SQL*Net, Distributed ORACLE relational

DBMS, and SQL*Connect.

SqL*Net provides distributed processing capability, and the application software can run on a different machine from the ORACLE kernel. Users can interface SQL*Net to arbitrary network environments. SQL*Net is communications protocol independent. It consists of a generic layer which is common to all protocols and environments, and a custom layer which is specific to each protocol and operating system.

The Distributed ORACLE relational DBMS supports location transparency, single site transactions, multi-site updates in multiple transactions and site autonomy. To meet System R* specifications, ORACLE is planning multi-site updates in a single transaction, replication and table partitioning for future releases. Distributed data dictionaries contain information about local data and how to get access to global data.

SQL*Connect allows ORACLE to communicate with non-ORACLE database management systems. Plans for broadening access to other database management systems, as well as to non-database file systems (VSAM, Record Management Service) are continuing. Query optimization is achieved by using an array interface to transfer data in batches. This insures that the network is used efficiently.

SALIENT FEATURES

ORACLE products run on an impressive range of mainframes, minis,

workstations and personal computers (over 90 hardware platforms), and support a large number of operating systems and a wide variety of network services. The designers intend to encrypt usernames, passwords, and SQL data. Approaches for meeting B1 and C2 National Computer Security Center (NCSC) security requirements are under development.

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2.19. Relational Technology, Inc.: INGRES

INGRES is marketed by Relational Technology, Inc. of Alameda, California. Products of this company are used extensively around the world. Approaches originally designed for homogeneous environments have been extended to cater to heterogeneous situations.

FUNCTIONAL ARCHITECTURE

The INGRES approach consists of: INGRES/NET Network Protocol Support, INGRES/STAR Distributed Data Manager, and INGRES Database Gateways. INGRES/NET is a transparent software layer that provides INGRES protocol support for DECnet, TCP/IP, SNA and ASYNC communications. This software allows access to databases across a network and performs automatic data conversion. With the INGRES/STAR Distributed Data Manager, simultaneous access to distributed databases across a network is possible. INGRES/STAR provides complete location transparency and local autonomy. INGRES Database Gateways enable other types of databases and file systems to be included in a distributed database. The appropriate gateway

module must reside on the local host. Relational, hierarchical, and network local data models are supported.

A two-phase commit protocol is used for distributed transaction processing. If a site involved in a transaction is inaccessible, other sites are automatically rolled back. When a system crashes, INGRES backs out all pending transactions and returns the database to a consistent state without operator intervention.

The global data manipulation language of INGRES is ANSI SQL. After receiving a SQL query from an application, INGRES/STAR breaks it into subqueries and directs them to the appropriate local data managers and gateways. Next, each local data manager executes its respective subquery and the data are transferred to INGRES. Data can be transparently joined and combined with INGRES SQL end-user interfaces or application development tools.

The INGRES/STAR Distributed Data Manager combines distributed query optimization with parallel processing. The distributed optimizer selects the best route for satisfying a multicomputer request. This approach reduces network traffic and improves performance. With this open, multi-server architecture, system administrators can customize their environment by assigning various servers with dissimilar priorities for different tasks, and by connecting transaction applications to a high priority server and batch report-oriented applications to a lower-priority server.

INGRES allows dynamic control of parameters to adjust the

concurrency of the system to the needs of the application set. Transaction and decision support applications may be executed at the same time. Also, the multi-server architecture allows systemwide monitoring and control of database sessions for real-time system management. A virtually infinite number of simultaneous users is possible.

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The INGRES data dictionary stores all performance related information-table structures and sizes, available indexes, and data distribution statistics. The more information is provided in the dictionary, the more intelligent is the Query Optimizer. The Query Optimizer produces query execution plans which can be examined to determine the effects of adding indexes, selecting different storage structures, and adding data distribution statistics.

SALIENT FEATURES

INGRES incorporates many of the features required for distributed transaction applications including two-phase commit, automatic recovery, and table fragmentation. Distributed query optimization strategies and full parallel query execution capabilities enhance its distributed on-line transaction processing capability.

2.20. Software AG: ADABAS

The ADABAS DBMS, developed by Software AG of West Germany, is a part of an open Integrated Software Architecture (ISA). Used in conjunction with NET-WORK (for distributed processing) and ADA-

NET (for distributed databases), ADABAS can support heterogeneous networks involving IBM, SIEMENS, DEC-VAX and WANG computers. The ADABAS DBMS supports relational, hierarchical, and network local data models.

ADABAS is designed for production, development, and information center environments. It claims to handle both high transaction volumes and complex ad hoc queries operating against large volumes of data. Distributed transactions can be in the form of a series of updates to different databases which might involve more than one network node.

FUNCTIONAL ARCHITECTURE

ADABAS uses either SQL embedded in COBOL, FORTRAN or PLI or NATURAL as its global data manipulation language. NATURAL is a fourth generation language based on a Communication Interface and standard protocols. NATURAL applications may be ported without modification across a wide range of system environments and hardware architectures.

The ADANET product of Software AG allows a user level logical ADABAS call to be translated into one or more physical calls in a transparent manner. A dynamically maintained ADANET translation table defines logical databases and their corresponding physical database representations. Three logical file structures are supported by ADANET: replicated, distributed and isolated files.

PREDICT is the centralized integrated active data dictionary. The structure of the metadata in PREDICT describes the logical and physical interconnections of all data and programs to be managed. PREDICT offers a user-friendly, menu-driven user interface for the management of data and application information.

SALIENT FEATURES

The Open Integrated Software Architecture (ISA) provides a high level of functionality in areas such as end user access, application solutions, application development and database management. It also provides software independence across multiple hardware and operating platforms without requiring any change in local databases. Global update processing is supported for relational, network and hierarchical local data models. The global data manipulation language of ADABAS reduces dependencies on third generation languages and is highly portable.

2.21. Sybase, Inc.: SYBASE

SYBASE is marketed by Sybase, Inc. of Emeryville, California. It can handle data in either hierarchical form or relational form. Data in a network DBMS are accessed via a relational interface.

FUNCTIONAL ARCHITECTURE

The Sybase Distributed System is comprised of three major components: the Sybase Data Server, the Sybase OpenServer, and

the Sybase StarServer. The Sybase DataServer allows remote queries and distributed transactions across multiple DataServers while the Sybase OpenServer extends this capability to other applications or DBMS Servers, including non-SQL systems. The Star Server adds the functionality of site transparency, including cross-site joins and views, and the Distributed Data Dictionary. The SYBASE OpenServer is in the final stage of beta testing; the DataServer and the StarServer are commercially available.

The SYBASE DataServer allows multi-site update transactions, while ensuring the logical consistency of the data via a two-phase commit protocol. A 'commit service' running at a designated site logs the progress of a transaction. The commit service logs notes about when all sites will be prepared. If the transaction fails before that point, it will be backed out; if it fails after that point, it will be rolled forward.

With its multithreaded client/server architecture, SYBASE assumes many of the functions of an operating system. SYBASE avoids the overhead of process swapping by scheduling and activating users. Multiple users share a single copy of the SYBASE image, data and procedure cache, and transaction log. Because each additional user requires only 30 KBytes of memory, hundreds of users can simultaneously access and update data on any node in a distributed network, according to Sybase.

A Virtual Server Architecture (VSA) for Symmetric Multiprocessor Systems is scheduled for beta testing. The VSA treats processors

as resources to be allocated on an as-needed basis. An application or a user is not restricted to a specific processor; instead applications may run on various processors at different points in time depending on the number of blocks for I/O requests or other reasons. Four key features of VSA are: synchronization control, automatic load balancing, optimal resource utilization, and dynamic configuration.

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SYBASE uses a Distributed Data Dictionary, consisting of synonyms which are used to locate objects, and Remote Object Catalogs which contain information about the characteristics of those objects. With this approach, users and applications need not be concerned with changes in the physical location of an object.

SALIENT FEATURES

Sybase is a leader in secure database management systems. The SYBASE Secure SQL Server meets Division B security requirements defined by the National Computer Security Center. In compliance with Division B specifications, SYBASE offers the ability to store data of multiple security classifications in a single database.

Because networking was built as a key part of SYBASE, rather than as an add-on, network usage is optimized. The network is used more efficiently by sending a single 'execute procedure' command, instead of an individual SQL statement.

SYBASE allows server-to-server communications. Multiple users may

share the communication link between servers, thus eliminating the overhead of opening and closing a linkage for each user.

2.22. Tandem Computers Incorporated: NonStop SQL

Tandem Computers Incorporated, headquartered in Cupertino, California, offers software and network compatible systems ranging from mainframes to low-cost network nodes. Because all Tandem processors use the same operating system and distributed database, as many as 255 systems with a maximum of 4,080 processors can be linked to form a geographically dispersed network. An individual system can support hundreds of terminals and a database with billions of bytes of data.

FUNCTIONAL ARCHITECTURE

GUARDIAN 90, a message based operating system, supports the parallel architecture of Tandem NonStop systems. Key features of GUARDIAN 90 are: fault tolerance, distributed data processing capabilities, high transaction throughput, linear expandability, security, data integrity, and the ability to link systems, resources, and devices.

GUARDIAN 90 supports fault tolerant process pairs. In the case where a system software error causes a process to fail, its backup process takes over in another processor. In addition, GUARDIAN 90 manages communication between processors. Database availability is always ensured. Mirrored volumes of data are stored on a pair

of physically independent disk drives. When a failed drive is operational again, the system automatically brings it up to date. The systems resume normal mirrored operation after the restored disk is fully updated. Via its message system, GUARDIAN 90 coordinates all distributed activity in a fashion that makes all data accesses transparent to the user. An application can read, write, and update data regardless of the physical location of the data. A built-in password security system is used to control file, device and process availability. Tandem offers integrated security and encryption products to protect information from unauthorized access.

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SALIENT FEATURES

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The multiple parallel processing capability, the fault tolerance capability, and the expandability feature make NonStop SQL a good product for on-line transaction processing (OLTP).

Tandem Corp. has signed integration agreements with Oracle Corp, Sybase, Inc., Relational Technology, Inc. (Ingres), and Information Builders, Inc. (FOCUS) to port their front-end tools to Tandem Systems. Details of the approaches used by these companies are presented in their respective subsections.

2.23. TRW: TDIE

TDIE is a data integration engine developed by TRW, Redondo Beach, California. Unlike conventional efforts which use a gateway with

query language mapping, TDIE has no global data manipulation language. A local DBMS is used and a connection to a remote DBMS is triggered when a user issues a demand requiring data integration.

TDIE is not oriented towards any particular environment. No change is required in existing applications. Additional local database management systems may be added to the system configuration and more than one global schema is permitted.

An integrated scheme performs uniform query and updates for hierarchical, relational and network data models. Automatic updating is performed by polling the potential source files at a rate suitable for a given data category. Types of data that are frequently accessed and must be current are polled at short intervals. This approach eliminates unnecessary activity and avoids unnecessary copying of bulk data by transferring only the changes in data.

FUNCTIONAL ARCHITECTURE

The TDIE is comprised of the Data Transform Manager and the Host Interface Manager. These may reside on one processor or they may be partitioned across multiple processors.

The Data Transform Manager contains the centralized dictionary which defines the global conceptual schema. This knowledge-based dictionary supplies scripts, translation software, and data

locations for full data integration. Scripts are instructions for data transfer, which are sent to the Host Interface Manager, while the translation software contains the translation of one query language to another.

To add another local DBMS, translators and scripts must be added to the dictionary. TRW provides a front end tool called the Integration Advisor for interactive construction of a global integrating schema.

The Host Interface Manager is the bridge to the host applications. It provides for automatic updating and data retrieval. Timecritical applications in transaction processing or real-time control are not supported. Also, because TDIE is not a distributed heterogeneous DBMS, concurrency control is an important issue.

SALIENT FEATURES

TDIE supports global updates and enables the integration of hierarchical, relational, and network data models. It is layered on top of the existing systems and does not require any change to them. New systems can be added as requirements change. Also, the TDIE can be implemented in increments as desired.

2.24. Xerox Advanced Information Technology: Multibase

Multibase was developed by Research and Systems Division of Computer Corporation of America. This division was acquired by

Xerox and is now called Xerox Advanced Information Technology (Xerox AIT). Although Multibase is a prototype and not a product, Xerox AIT utilizes this technology for providing customized distributed heterogeneous database management systems for its clients.

Multibase provides a uniform, integrated interface for retrieving data from several existing database management systems. It allows a user to reference data in heterogeneous databases, through a common query language, using a single conceptual schema. Multibase is designed to serve as a general tool, without specific orientation towards any particular application area. It allows existing applications to operate without change and also permits new local systems to be included in an existing multibase system configuration.

The integrated access available through Multibase does not provide either the capability to update data in local databases or the ability to synchronize read operations across several sites. In order to process user queries, the system must request and control specific services offered by the local systems (e.g., locking local items).

FUNCTIONAL ARCHITECTURE

Multibase uses the language DAPLEX as its global data manipulation language. DAPLEX provides constructs that allow users to model real world situations in an efficient manner. A three level schema

of definitions is employed as discussed in Section 1.5. The process of global information retrieval involves two main components. These are: Global Data Manager (GDM), and Local Database Interface (LDI). The GDM performs tasks of Command Processor, Decomposer, Merger and Distributed Execution Monitor. The LDI design of Multibase is based on the needs of a local DBMS, for example, a more sophisticated LDI is designed to support a file system as compared to the one that supports a DBMS.

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SALIENT FEATURES

Multibase provides an integrated scheme for uniform query access to dissimilar systems including hierarchical, relational, and network local data models, and File Systems. Global query optimization is performed in Multibase.

SECTION THREE: DISCUSSION AND CONCLUSION

The first section of this report explained the basic principles of integrating disparate information systems. In the next section, a broad spectrum of commercially-marketed products and approaches were discussed. Based on the material covered so far, this section develops a set of recommendations for selecting the optimal product or products. In particular, the information presented in Section 2 and Appendix B is utilized to develop a taxonomy for classifying all the alternative approaches.

In Section 1.4, it was emphasized that changes to existing information systems should be as little as possible; in the ideal case, no change should be required. Subsequently, in Section 1.5, the importance of a single global conceptual schema was emphasized. Many products and approaches tabulated in Appendix B indicate that multiple global schemata can be created to suit the needs of From a technical viewpoint, these are external various users. schemata, not global schemata. These are some of the considerations that were examined while developing the broad categories described in the succeeding subsections.

3.1. Unintrusive Integration

This is the most elegant form of integrating information systems. The word "unintrusive" connotes that the integration methodology does not demand that specific modules be resident in each

participating information system (see Figure 1). There is a definite assumption in products belonging to this category that existing information systems cannot, and should not, be tampered with. On top of these systems, a single global conceptual schema is introduced.

III

Use of this approach in a multiorganization or a multidivision environment implies that the existing information systems of various divisions and organizations will not require any modification. These systems will be accessed by a central system that holds a global schema with names and locations of all data items in participating systems. A user will access the central system, which will retrieve information from appropriate sources. The use of a central system and a single global schema offers appropriate controls on data integrity, access rights, and the feasibility of developing a new information system over time.

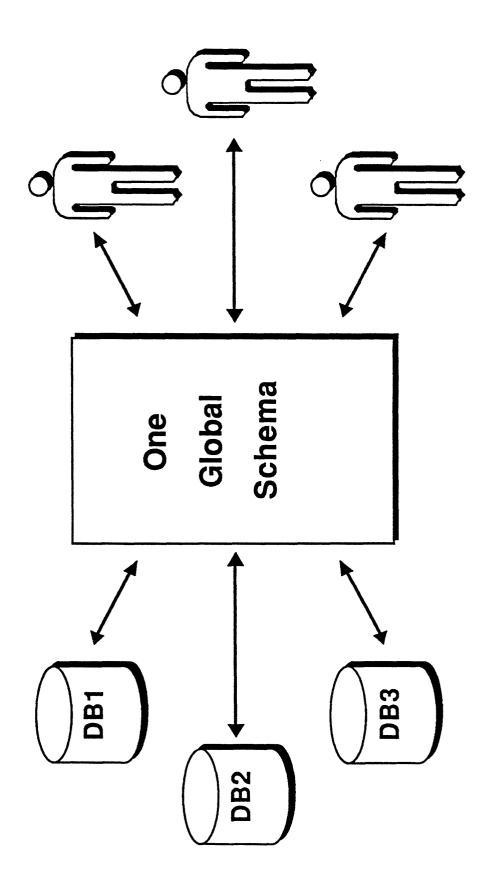
The main disadvantage of this approach is that is difficult to make global updates to multiple systems as there is no control over the capabilities of individual information systems. The application program must, therefore, be designed to safeguard the integrity of the updates.

Notable products in this category are: Multibase (developed by Computer Corporation of America and now marketed by Xerox); Mermaid (developed by Unisys and now marketed by Data Integration, Inc.); and CDM*Plus (developed under the IISS initiative of U.S.

Unintrusive Integration

Figure 1:

(No change to local information systems)



Air Force and now marketed by Control Data Corporation). The development of all these three products involved huge sums of money, and the acceptance of these products is still not commensurate with the investment. Defense organizations have been in the forefront in evaluating and field-testing these products.

3.2. Intrusive Integration

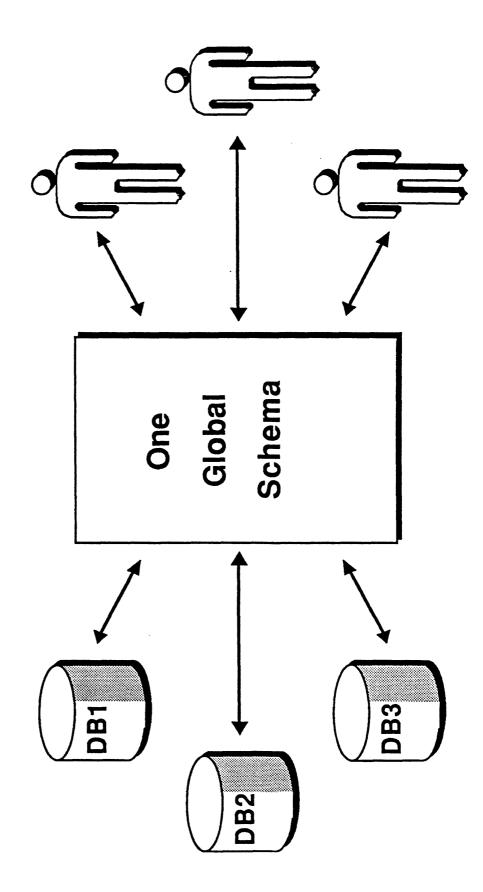
Approaches in this category require some modification of existing information systems or the location of specific modules within these systems. These modules are typically utilized to perform network services, to store information on location of specific files, or to initiate secondary processes in case some information resources are not operational or unusually slow in responding to requests (see Figure 2). If these modules are designed and implemented properly, and if the existing information systems offer the flexibility to accept these modules from technical and organizational considerations, then this approach can offer significant performance improvements over the approach outlined in Section 3.1. The main drawback, however, lies in the fact that existing information systems frequently use older technologies, employ outdated programming practices, offer little good documentation, and can only be modified after prolonged effort.

Notable products in this category include SUPRA (original technology developed in West Germany and now marketed by CINCOM), INGRES (from Relational Technology Inc.), and Sybase (from Sybase).

Intrusive Integration

Figure 2:

(Some changes to local information systems)



In the case of Sybase, it appears that the present family of products support SQL databases only, and support for non-SQL based databases and file systems is still being tested. There are some restrictions on the types of data models and data formats supported by these three products. For example, Supra requires data to be in first normalized form and keys to be in ascending or descending order. CINCOM is interested in developing a common protocol for integrating heterogeneous databases, that could be utilized by multiple companies. Because of these reasons, it is important to analyze potential applications carefully to determine which of these products, if any, is appropriate in the short-term and in the long-term.

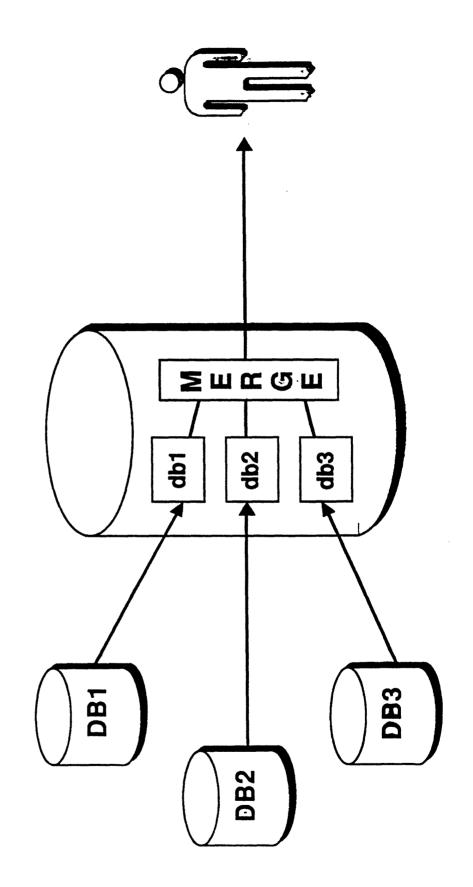
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3.3. Interfacing with Controlled Redundancy

In this approach, no effort is made to create a global conceptual schema. Instead, based on the perceived needs of the user, strategies are designed to get the necessary pieces of information on a "just-in-time" basis (see Figure 3). As an example, consider a hypothetical application which requires that information on 100,000 items (which could be key customers, spare parts, or inventory levels) be available every morning. Assume that the existing information systems that are used to track the position of these items are updated with different periodicity: very few on a yearly basis, some on a monthly basis, a large percentage on a weekly basis, another large percentage on a daily basis, and the rest on a six-hourly basis. In such a situation, information can



Interfacing with Controlled Redundancy (Selective "Download" and Integrate)



be obtained in advance; further, to fulfil anticipated requests one can get information only for the items that have changed since the previous status report.

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One product that exemplifies this approach is Data Integration Engine from TRW. In this patented approach, when a change in the status of a source information system occurs, this change is transmitted to the destination application with a priority that would allow it to reach the destination before the next scheduled time of retrieval. The software to accomplish this task and other related tasks is comprised of two types of material: a basic data engine, which is site independent, and additional material which must be customized to each site. The basic data engine is written in C and can operate in a UNIX or a VAX environment. The additional material can be generated with an expert system called Integration Advisor, which was originally written in LISP and has recently been converted to C. According to a technical representative of TRW, the ratio of the size of the basic data engine to the size of the additional material is 5:1. Incidentally, TRW has recently been awarded a multiyear contract by the U.S. Department of Defense.

3.4. Interfacing with Transaction Transfer

This is the simplest technique for managing transactions in a heterogeneous information system environment. Instead of creating a global schema or prestaging data to meet expected needs, this

approach implies that action begins after a user enters a command. Unless the user is willing to tolerate long delays, this approach is good only in situations where only a small number of information systems need to be accessed in response to the user query (see Figure 4).

Among the notable products in this category, the one that deserves mention here is DIS from Metaphor Computer Systems, Inc. DIS possesses two unique characteristics. First, it is completely icon-driven and as such it is very easy to learn and use. Second, IBM has invested in Metaphor Computer Systems Inc. and they are working together to provide a similar interface for IBM mainframes. As such, the DIS interface has the potential of becoming a commonly used standard for dealing with multiple information systems.

3.5. Additional Evaluation Criteria

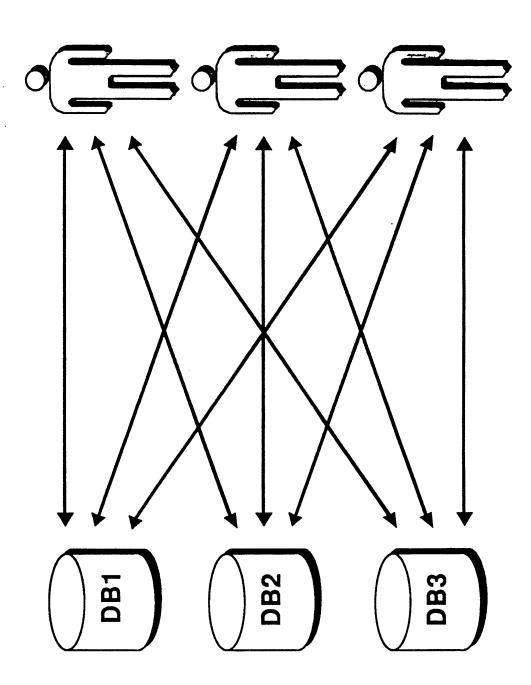
In addition to determining which of the four distinct approaches described in Sections 3.1-3.4 above may be most appropriate for a particular application, it is necessary to examine a number of other technical and non-technical issues before embarking on a specific implementation strategy. These issues are delineated in the following paragraphs.

3.5.1. Composition of Transactions

Transactions are of two types: queries and updates. In our



Interfacing with Transaction Transfer



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context, a query involves reading of data in multiple computerbased systems, manipulating the data retrieved to make them consistent, and then presenting the result to the user. The original data in source computer-based systems remain unaltered by a query. An update, however, results in modification of data in various information systems. It is much more difficult to deal with updates than with queries. If an application involves updates, then a significant number of the alternative approaches discussed in Section 2 should be deleted from consideration. The ones to be deleted can be readily identified from Appendix B, which contains a specific question on how update transactions are managed.

3.5.2. Security Considerations

When one makes it easier to automatically access relevant pieces of information in disparate systems, the same capabilities become available to any unauthorized user of the system. Any integrated set of systems is more vulnerable to security violations than a single information system.

None of the vendors was specifically asked to comment on this issue, although a few of them have provided information on their compliance level with security requirements defined by the National Computer Security Center. This issue may hold significant weightage in the case of large defense, banking, and insurance applications.

3.5.3. Corporate Expertise

Most large computer corporations have tended to concentrate on integrating systems of their respective makes only. For example, in Section 2, there was a discussion of IBM and DEC suggesting solutions pertaining to their respective families of products. Frequently, smaller companies come up with non-traditional concepts, and larger companies tend to endorse them after these concepts have attained some degree of maturity and acceptance by the market.

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The fact that the companies who developed good prototypes such as Multibase and Mermaid later spun off their development groups must be viewed with caution. This is an evolving field, and user organizations must be careful about selecting only the set of products and services that are likely to be supported by more than one company.

3.5.4. Applications and Development Strategy

In the current business environment of increasing globalization and corporate reorganizations, applications and functions are apt to change in relatively short periods of time. The present technical report has been written without any specific application in view. As an organization determines which applications are important in the short-term and the long-term, it would become feasible to make more specific comments about the relevance of

different products and approaches.

It is important to note that any endeavor towards integration has ramifications for future integration initiatives. The selected approach must offer the flexibility to be enhanced in a modular fashion. Further, public-domain approaches should be preferred over proprietary ones. Based on these considerations, a development strategy can be delineated.

3.6. Conclusion

The integration of existing information systems involves a careful analysis of many technical, organizational, and strategic issues, and the surmounting of many technical and nontechnical barriers. This technical report has focused on products and their capabilities related to logical access to relevant pieces of information. For discussion of other approaches and issues, the reader may wish to refer to Ref. [2] and the set of eight books on Knowledge Based Integrated Information Systems Engineering published by M.I.T.

The problem of heterogeneity can be overcome either by integration or by interfacing. To decide which option to select, it is worthwhile to look back to the stage when the concept of database management system originated. If one wanted to deal with a single clearly-defined application, it was better to write a customized program to do so; however, if there were many users with diverse

needs, it became increasingly attractive to use the concept of a database management system, even though it involved a larger investment at the initial stage. The same considerations apply at the metalevel today. If access to multiple information systems is needed exclusively for a small set of clearly-defined applications, the interfacing alternative should suffice. However, if there is likely to be an array of broad applications, it appears preferable to take the integration alternative. Remember that "just as the evolution of integrated database management technology displaced conventional piecemeal programming approaches in the 1960s, similarly the new information integration alproaches will lead to a significantly different computational environment by the year 2000 A.D.!" [1].

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Appendix A

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COMPANIES CONTACTED AND APPLICABILITY

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APPLICABILITY	Relevant. Upon completion of contract negotiations, AT&T will provide more in- formation.	Not relevant. ACE is an educational institution with no product or R&D.	Not relevant. RDM/2 is a relational DBMS. Integra- tion of other distributed databases is not currently within its capabilities.	Relevant. Details includ- ed in document.
PRODUCT/ SYSTEM	Rapid	N/A	RDM/2	Cronus Distributed Computing Environment
CONTACT PHONE & ADDRESS	Doug Justice Marketing Interstate 85 at Mount Hope Church Road McLeansville, NC 27301 Phone: (919) 279-3491 FAX: (919) 279-3980	Beverly Phone: (415) 941-3399	Nardis Martinez RDB Product Manager 9232 Eton Avenue Chatsworth, CA 91311 Phone: (818) 998-7666	Steve Vinter Cronus Product Manager 10 Moulton St. Cambridge, MA 02138 Phone: (617) 873-2876 FAX: (617) 873-3776
COMPANY	АТ&Т	Advanced Computing Environments	Amperif Corporation	BBN Systems and Technologies Corporation

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<u>APPENDIX A</u> COMPANIES CONTACTED & APPLICABILITY

COMPANY	CONTACT PHONE & ADDRESS	PRODUCT/ SYSTEM	APPLICABILITY
Booz Allen & Hamilton, Inc.	David Subar Senior Consultant Artificial Intelligence Practice 4330 East West Highway Bethesda, MD 20814 Phone: (301) 951-4580	SMIDI	Relevant. Details includ- ed in document.
CRI, INC.	No Contact Phone: (408) 980-7499	Relate/DB	Unknown. Answering Ser- vice takes calls but none returned.
Cambridge Technology Group	Steven Schenefeld 219 Vassar Street Cambridge, MA 02139 Phone: (617) 876-2338 FAX: (617) 499-1777	Surround	Relevant. Details includ- ed in document.
Cincom	Ken Pagnotta Regional Technical Manager 35 Braintree Hill Park Suite 100 Braintree, MA 02184 Phone: (617) 849-1020 FAX: (617) 849-1030	Supra	Relevant. Details includ- ed in document.

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 APPENDIX
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 COMPANIES
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COMPANY	CONTACT PHONE & ADDRESS	PRODUCT/ SYSTEM	APPLICABILITY
Cognos Corporation	Judd Lowe Senior Sales Representative 2 Corporate Place I-95 Peabody, MA 01960 Phone: (508) 535-7350 FAX: (508) 535-8241	PowerHouse	Relevant. Details includ- ed in document.
Computer Associates International Inc.	Jack Hadder District Manager, Defense 8227 Old Court House Road Vienna, VA 27181 Phone: (703) 356-7700 x635 FAX: (703) 847-0290	Datacom/DB	Not relevant. Datacom/DB does not support hetero- geneous database environ- ments.
Computer Corporation of America	Pete Goldman Senior Technical Representative 1800 Diagonal Road Alexandria, VA 22314 Phone: (703) 836-5200 FAX: (703) 836-6590	Multibase	The Research and Systems Division which developed Multibase has been sold to Xerox. See appropriate entry below.
Consultants for Management Decisions, Inc.	Todd Collins One Broadway Cambridge, MA 02142 Phone: (617) 497-4434 FAX: (617) 497-2305	N/A	Not relevant. Company builds custom software systems to serve complex information needs of large organizations.

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COMPANY	CONTACT PHONE & ADDRESS	PRODUCT/ System	APPLICABILITY
Control Data Corporation	William Bradley CDM plus Business Development Integration Technology Services 2970 Presidential Dr., Suite 200 Fairborn, Ohio 45324 Phone: (513) 427-6323 FAX: (513) 427-6301	CDM ^t plus	Relevant. Details includ- ded in document.
Control Data Corporation	Roger West Seminary Rd. Suite 1000 Alexandria, VA 22311 Phone: (703) 998-3575 FAX: (703) 998-3571	Ascent (Intelligent Gateways)	Relevant. Details includ- ed in document.
Cullinet	Alan Miller 400 Blue Hill Drive Westwood, MA 02090 Phone: (800) 551-4555	IDMSR & Enterprise Database	Not relevant. IDMSR runs only on IBM 370 Architec- ture and Enterprise Data- base runs only on the VAX.
DB/Access, Inc.	Mike Bergeron Director OEM Sales 2900 Bordon Ave. Suite 101 Santa Clara, CA 95051 Phone: (408) 735-7545 FAX: (408) 735-0328	Access Star	Not relevant. Access Star does not support data in- tegration.

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COMPANY	CONTACT PHONE & ADDRESS	PRODUCT/ System	APPLICABILITY
Data Access Corporation	Katherine L. Cronin Manager Government Accounts 1400 S.W. 119th Ave. Miami, FL 33186 Phone (305) 238-0012	Dataflex	Not relevant. The product is a relational DBMS, but does not provide the abil- ity to integrate distrib- uted databases.
Data Integration, Inc.	Marjorie Templeton President 3233 Federal Avenue Los Angeles, CA 90066 Phone: (213) 390-4507	Mermaid	Relevant. Details includ- ed in document.
Digital Equipment Corporation	Glen Light Government Software Services 721 Emerson Road St. Louis, MO 63141 Phone: (314) 991-6552 FAX: (301) 731-3018	Rdb/VMS, DBMS	Relevant. Details includ- ed in document.
Electronic Data Systems	Craig Rogers 13600 EDS Drive Herndon, VA 22071 Phone: (703) 742-1572 Phone: (703) 742-1555 FAX: (703) 742-1325	Custom- designed applications	Initial review suggested further investigation. The contact was requested to fill in a Standard Technical Specification Sheet on 8/29/89.

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APPENDIX A COMPANIES CONTACTED & APPLICABILITY

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COMPANIES	
CONTACTED	APPENDIX
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COMPANY	CONTACT PHONE & ADDRESS	PRODUCT/ System	APPLICABILITY
General Motors Research Laboratories	Chin-Wan Chung GM Research Laboratories Warren, MI 48090 Phone: (313) 986-1386	Dataplex	Unknown. GM is working with EDS and several DBMS vendors on a commercial product. This product is still in the prototype stage.
Gupta Technologies	John Ambler 12200 Sunrise Valley Drive Reston, VA 22091 Phone: (703) 264-9002 FAX: (703) 264-0568	The SQL System	Relevant. Details includ- ed in document.
Hewlett Packard	John Blue 1550 Wilson Blvd. Arlington, VA 22209 Phone: (703) 284-2635	Custom- designed applications	Initial review suggested further investigation. The contact was requested to fill in a Standard Technical Specification Sheet on 8/17/89.
Honeywell Federal Systems, Inc.	Hal Hannickel 7900 West Park Drive Mail Stop 603 McLean, VA 22102 Phone: (703) 827-3776 FAX: (703) 827-3729	DDTS	Relevant. Details includ- ed in document.

Company	CONTACT PHONE & ADDRESS	PRODUCT/ SYSTEM	APPLICABILITY
IBM	Barbara Lallas IBM / 5C1 1 Copley Place Boston, MA 02116 Phone: (617) 638-2770 FAX: (617) 638-1735	SAA, DB2, and SQL/DS.	Relevant. Details includ- ed in document.
Information Builders, Inc.	Michael Howard FOCNET Product Manager 1250 Broadway New York, NY 10001 Phone: (212) 736-4433 FAX: (212) 967-6406	FOCNET	Relevant. Details includ- ed in document.
Information Dimensions, Inc.	Debbie Fitzgibbons 50 Stanford St. Suite 800 Boston, MA 02114 Phone: (617) 227-7155	DM	Not relevant. The product is a relational DBMS, which integrates homoge- neous, relational database management systems.
Informix Software, Inc.	Donna Brewer Regional Sales Representative 4100 Bohannon Drive Menlo, Park, CA 94025 Phone: (415) 322-5166	Informix SQL	Not relevant. The product is a relational DBMS. Distributed capabilities will not be available un- til September 1990.

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APPENDIX A COMPANIES CONTACTED & APPLICABILITY

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COMPANY	CONTACT PHONE & ADDRESS	PRODUCT/ SYSTEM	APPLICABILITY
Interbase Software Corporation	Brian Smith Senior Sales Representative 209 Burlington Rd. Bedford, MA 01730 Phone: (617) 275-3222	Interbase	Relevant. Details includ- ed in document.
Intercim	Jim Padden Product Manager 4030-2 West Braker Lane Suite 200 Austin, Texas 78759 Phone: (512) 345-8872 FAX: (512) 345-0305	FDL/FDM	Not relevant. Product in- tegrates manufacturing ap- plications to one rela- tional DBMS to provide factory decision control systems.
Intersystems Corporation	Steve Reagon Sales Representative 1 Memorial Drive Cambridge, MA 02139 Phone: (617) 621-0600	M/SQL	Not relevant. The product is a relational DBMS but integrates homogeneous relational database management systems only.
Massachusetts Institute of Technology	Richard Wang Assistant Professor Sloan School of Management Phone: (617) 253-0442	CIS	Relevant. Details includ- ed in document.

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COMPANIES CONTACTED	APPENDIX A
& APPLICABILITY	A

COMPANY	CONTACT PHONE & ADDRESS	PRODUCT/ System	APPLICABILITY
Metaphor Computer Systems	Ron Goldstein 2010 Corporate Ridge Suite 510 McClean, VA 22102 Phone: (703) 749-3620 FAX: (703) 749-3684	DIS	Relevant. Details includ- ed in document.
Must Software International	David Chester 1030 Massachusetts Ave. Cambridge, MA 02138 Phone: (617) 868-2950 FAX: (617) 876-7918	Nomad	Initial review suggested further investigation. The contact was requested to fill in a Standard Technical Specification Sheet on 8/14/89.
On-Line Software International	Peter MacPherson 75 Second Avenue Needham, MA 02194 Phone: (617) 449-8860	Ram's Information System	Relevant. Details includ- ed in document.
Oracle Corporation	Christie Ellsworth Technical Account Manager 3 Bethesda Metro Center Suite 1400 Bethesda, MD 20814 Phone: (301) 907-2739	ORACLE	Relevant. Details includ- ed in document.

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COMPANY	CONTACT PHONE & ADDRESS	PRODUCT/ SYSTEM	APPLICABILITY
Plus 5 Computer Services	Carter Crandal 1968 Innerbelt Business Center Drive St. Louis, MO 63114 Phone: (314) 426-3900	Unix Mumps	Not relevant. The product is not a DBMS, does not integrate or translate data, and provides no global schema.
RIM Technology, Inc.	Bill Holter Marketing Manager 1775 12th Avenue NW Suite 115 Issaquah, WA 98027 Phone: (206) 392-4776	RTRIM	Not relevant. The product is a DBMS limited to inte- gration of homogeneous, relational databases.
Relational Technology, Inc.	Lori Dreyfus 6 New England Executive Park Burlington, MA 01803 Phone: (617) 727-5060 FAX: (617) 272-3135	Ingres Star	Relevant. Details includ- ed in document.
Sharebase (Formerly Britton Lee)	Paul Bass Director, Technical Marketing 14600 Winchester Blvd. Los Gatos, CA 95030 Phone: (408) 378-7575 FAX: (408) 866-0774	Sharebase	Not relevant. The product integrates homogeneous database management sys- tems only.

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COMPANY	CONTACT PHONE & ADDRESS	PRODUCT/ SYSTEM	APPLICABILITY
Smith, Abbott & Co.	Ed Geier 1725 K Street NW Suite 501 Washington D.C. 20006 Phone: (202) 223-4990	Autopro	Not relevant. The product performs distributed management only for its own databases.
Software AG of North America, Inc.	Patti Malanka Federal Systems, Inc. Suite 210 11490 Commerce Park Dr. Reston, VA 22091 Phone: (703) 620-0100	ADABAS	Relevant. Details includ- ed in document.
Sybase	Peter Loux Federal Programs Manager 4600 South Ulster St. Suite 700 Denver, CO 80237 Phone: (303) 721-3308	Sybase	Relevant. Details includ- ed in document.
TRW	Tony Materna Systems Engineering & Devel. Div. One Space Park Bldg. DH2/Room 2328 Redondo Beach, CA 90278 Phone (213) 812-6021	TDIE	Relevant. Details includ- ed in document.

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APPENDIX A COMPANIES CONTACTED & APPLICABILITY

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Appendix B

DETAILED CHARACTERISTICS OF PRODUCTS AND APPROACHES

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<u>APPENDIX A</u> COMPANIES CONTACTED & APPLICABILITY

COMPANY	CONTACT PHONE & ADDRESS	PRODUCT/ SYSTEM	APPLICABILITY
Tandem Computers, Inc.	Dave Magram Mail Stop 219-44 5300 Stevens Creek Blvd. San Jose, CA 95129 Phone: (408) 553-6757 FAX: (408) 553-6960	NonStop SQL	Relevant. Details includ- ed in document.
Unify Corporation	Frank Verardi 3870 Rosin Court Suite 100 Sacremento, CA 95834 Phone: (916) 920-9092	Accel1/SQL	Not relevant. The product is a DBMS with a 4GL lan- guage, but does not in- tegrate heterogeneous database management sys- tems.
Unisys	Anita Holmgren Deputy Dir. of Advanced Projects Network and Information Securi- ties Division 5151 Camino Ruiz Camarillo, CA 93010 Phone: (805) 987-9300	Mermaid	Rights to Mermaid have been licensed to Data In- tegration, Inc. See appropriate entry above.
Xerox Corporation	Gene Kakalec Adv. Information Technology Dept. 1800 Diagonal Rd. Suite 300 Alexandria, VA 22314 Phone: (703) 836-1823	Multibase	Relevant. Details includ- ed in document.

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-	Global Data Manipulation Language	Local DBMS Supported	Local Data Models Supported	Environment (Software and Hardware)	Acronym Derived Prom	
	SQL Subset	Proprietary object database; possible to integrate other databases; Informix and Oracle are currently supported; ven- dor has experience with Ingres and Vbase.	Object-oriented and Relational	General. SUN 2, 3, 4, 386i: SUNOS, and Mach 2.0 (on SUN 3). VAX Pamily: Ultrir, VMS, Masscomp RTU, Symbolics Genera, C, Portran and CommonLisp are supported.	Моле	BBN SYSTEMS & TECENOLOGIES: Cronds
	The Query Processor supports a high level, iconic, global query language. Pacilities can be added to support natural language, if appropriate.	See remarks.	Hierarchical, Relational, Net- work, and Object Oriented models can be supported through the Database Models Knowledge Base.	General.	Intelligent Distributed Information Management System	BOOJ ALLEN: IDINS
	Any user interface, icons, voice, SQL, graphics, forms, etc.	IMS, IDMS, Oracle, DB2, and all others on remote systems. Ingres, Oracle, Informir, Sy- base and all others on local Surround bor.	All, including flat file, hier- archical, relational and net- work, via Surround SM interface and coordination tools.	General. The Surround tools have been ported to virtually every UNIX platform including NCR, HP, SUW, IBM, AT&T, Unisys, and Digital for interconnectivity and added functionality.	Wone	CAMBRIDGE TECHNOLOGY GROUP: SURROUND
-	SõF	SUPRA DPDM and VSAM.	- Relational and Network.	General. DEC, IBM, WIXDORP	None	CINCON: Supra

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	COGNOS CORPORATION: POWERHOUSE	CONTROL DATA CORP: ASCENT*suite	CONTROL DATA CORP: CDM*PLUS	DATA INTEGRATION, INC.: MERMAID
	None	None	Common Data Model * Plus	Иове
Acronym Derived Prom				
Environment (Software and Hardware)	General.General.PowerHouse 4GL on DEC VAX,A UNIX-based software productHP3000, HP9000, DG MV (AOS) and ported to a number of hardwareIBM PC and PS/2 (DOS and OS/2).PowerHouse 4GL on DEC VAX,	General. A UNIX-based software product ported to a number of hardware platforms including Pyramid,	General. CDM*plus operates on VAX/VMS, HP/UX, Pyramid/OSr. CDM*plus accesses applications and data-	General. Portable - written in 'C'. Server with optimizer/control- ler/network on a UNIX processor
and Bardware		prations including Fyramid, Gould, Cubix, DEC, WCR Tower, ARIX, ELisi, CDC910, SUN, and Celerity.	<u>ہ</u>	apprications and data- retinetwork on a unit processor the above systems plus such as SUN. User interface on nd Amdahl/MVS. UNIX system available. DBMS on diverse hardware and OS.
Local Data Models Supported	Sequential, Relative, and In- dered files; Network and Rela- tional databases.	This depends on the UNIX data- base supported by the UNIX platform being run on.	Relational, Wetwork/Codasyl, and Hierarchical.	Relational and Network. Also, supports objects in files. Qualification to locate files is done on a structured rela- tion that contains the file name. Retrieve process and display process are specific to object type.
Local DBMS Supported	Image3000, TurboImage, KSAM, MPE on HP. C-ISAM on HP UX. RMS, Rdb/VMS, PowerHouse Star- Base on VAX/VMS. INFOS, ISAM, INFOS DBAM, DG/SQL on DG AOS. PHISAM C-tree Deriv. on IBM PC.	See above.	Oracle, Ingres, DB2, IDMS on IBM/MVS. Oracle, Ingres, VAX DBMS on VAY/VMS. Oracle on Pyramid. Oracle on HP. Others on request.	IDM (Sharebase) on VAX. Ingres on SUN 120, SUN 170 or other UNIX.
Global Data Manipulation Language	PowerHouse 4GL; high-level op- erations such as join (inner and outer), restriction (selec- tion) and sorting. Powerhouse StarBase: fully relational DBMS supporting SQL and GROTON DML.	Would have to be programmed using C, and Shell Scripts in combination with *gateway scripts.	Neutral Data Manipulation Lan- guage (NDML), a SQL based neu- tral query language. NDML can be embedded in Cobol, Portran, and C.	1ðS

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	DIGITAL EQUIPMENT CORP.: Rdb/VMS, DBMS	GUPTA TECENOLOGIES, INC: SQL SYSTEN	HONETWELL: DDTS	IBM: SAA, DB2, AND SQL/DS
Acronym Derived Prom	None	Яоде	Distributed Database Testbed System	DB2 - DATABASE 2. SQL/DS - Structured Query Lan- guage/Data System
Environment (Software and Hardware)	General. VAX/VMS.	General. DOS 3.1 through 4.1, OS/2, and UNIX on a SUM platform.	General. Bull-HN's DPS/G and DPS/88; IBM-PC; Charles River's Mork- station MOD-400 and GCOS8; DOS; UNIX	Genera].
Local Data Models Supported	Hierarchical/Wetwork and Rela- tional	Relational (ertended). SQL Level I + enhancements are sup- ported.	Wetwork (CODASYL) and Rela- tional.	Relational and Hierarchical.
Local DBMS Supported	DBMS and Rdb/VMS.	SQLBase is both a local and a client-server database. SQL- Base is compatible with IBM's DB2 and SQL/DS. Connections for Oracle, DB2, and SQL/DS are performed by Gupta tools.	IDS-II, MDBS-III, and RAM (Re- lational Access Manager).	SQL/DS on IBM/VM. MVS.
Global Data Manipulation Language	SQL, RDO, DSRI-Rdb, DBMS-MML.	SQL with enhancements (SQL- Talk), pre-compiler for ANSI COBOL, and application program interfaces for C.	GORDAS and SQL	1õS

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Global Data Manipulation Language	Local DBMS Supported	Local Data Models Supported	Environment (Software and Hardware)	Acronym Derived Prom	
3GL, 4GL FOCUS language, and ANSI SQL.	ADABAS,C-ISAM,DATACOM/DB,DBMS, DB2,Dbase,DMS,DMS/TX,MCBA,PACE, SPEED2,FOCUS,IDMS,IDMS/R,In- formix,INGRES,ISAM,Model 204, Oracle,QSAM,System 2000,Rdb, RMS,Sharebase, and others.	Hierarchical, Relational, and Network.	General. IBM PC/DOS, IBM/VM, IBM/MVS, VAX/VMS, WANG/VS, MPE/X, MPE XL (HP), Guardian 90, UNIX, various other platforms and operating systems.	POCUS Network Products.	INPORMATION BUILDERS, INC.: POCNET
BLR, a lower level language based on the Digital Standard Relational Interface is a mes- sage oriented language designed to be generated and read by programs.	Interbase on VAX, Apollo, SUW, HP, OS/2 Rdb/VMS.	Relational.	General. Distributed minicomputer and workstations. Networked applications in UNIX and VMS environments. Compatible with VAX/VMS, SUN, Apollo, HP, and OS/2.	Interconnective, heterogeneous hardware, operating system and database.	INTERBASE SOPTWARE CORP.: INTERBASE
Entity-Relation (ER) - like query language.	Informix-SQL on AT&T 3B2.VAX (VMS) - Oracle, Ingres,Oracle-SQL on IBM PC/RT. SQL/DSSybase, and Rdb, IBM - SQL/DSon IBM 4381.Finsbury Datalineand DB2, Teradata DBC, Share-(menu-driven Hierarchical)Base 700 and Metaphor DBS.I.D. Sharp Disclosure (propri-Iconic tools to merge data froetary command-driven language).PC Dif, ASCII and Lotus files.	Relational, Menu-driven, and Relational and relational Hierarchical. Conversion between a local database and non-relational DBMS facili CIS is done through a Local through TTY, 3270, and 378 Query Processor (LQP) operating emulators that are invoked on the CIS processor. click actions by the user.	General. Written in LISP, C, and UNIX shell. Runs on UNIX System V. DBMS on diverse hardware and software.	Composite Information Systems.	MASSACHUSETTS INSTITUTE OF TECHNOLOGY: CIS
User points-and-clicks at same iconic interface across all supported relational DBMS. SQL automatically produced. Meta- phor gateway servers convert SQL as required for platform.	-SQL on AT&T 3B2. VAX (VMS) - Oracle, Ingres, 2L on IBM PC/RT. SQL/DS Sybase, and Rdb, IBM - SQL/DS 181. Finsbury Dataline and DB2, Teradata DBC, Share- iven Hierarchical) Base 700 and Metaphor DBS. 2P Disclosure (propri- 1 Disclosure (propri- maand-driven language). PC Dif, ASCII and Lotus files.	Relational and relational in- terface. Interface to other non-relational DBMS facilitated through TTY, 3270, and 3780 emulators that are invoked through simple mouse point-and- click actions by the user.	General. The DIS workstation can be a combination of PC/XT, AT, or 386s running DOS 3.0 or higher, most PS/2 models running DOS, and PS/2 M70 & 80 running DOS or OS2.	Data Interpretation System.	METAPHOR DIS

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1867	2/87	"The Economics of Software Quality Assurance: A System Dynamics Based Simulation Approach," Abdel-Hamid, T., and Madnick, S.
1864	2/87	"A Temporal and Spatial Locality Theory for Characterizing Very Large Data Bases," Madnick, S., and Moulton, A.
1861	2/87	"Key Concerns of Executives Making I/S Decisions," Madnick, S., and Wang, Y.R.

Please Use Attached Form When Ordering Papers

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References	Remarks	
Purnished upon request.	The RAMIS Information System Product includes a DBMS, a Data Maintenance Facility, a Report Writer, a compiled ap- plication development lan- guage, and a developers tool- kit. The developers toolkit provides dictionary utilities, application creation, and re- lational capabilities. Also, a full function 4GL/DBMS for the personal computer with Mi- cro-Mainframe communications is available.	ON-LINE SOPTWARE INTERNATIONAL: RAMIS INFORMATION SYSTEM
Purnished upon request.	Oracle provides true hetero- geneous, distributed processing. mitted to deliver the best product is designed to give developers fleribility to build relational database system on cations easily and quickly. Additional tools are available for the end-user who is not equipped or interested in learning the SQL*Language.	ORACLE CORP.: ORACLE
Purnished upon request.		RELATIONAL TECHNOLOGY: INGRES
<pre>Pederal Aviation Admin. Mr. Rod Phillips (405) 680-3422 U.S. House of Representatives Mr. Jim Daley (202) 226-6426 Center for Disease Control Ms. Carolyn Seagraves (404) 639-0475.</pre>	The Software AG Product line allows for creation of truly distributed data processing systems with full user trans- parency.	SO ptware Ag: Adabas

References	Remarks	
Air Porce Military Airlift Com- mand Air Porce Strategic Air Command Air Porce Intelligence Agency	Open server designed to provide As many as 255 systems (4,080 users with toolset for building processors maximum) can be custom servers to heterogeneous linked to form a geographical databases. databases.	SYBASE: SYBASE
Tandem Reference Material U.S. Pat - Introduction to WonStop SQL awarded - WonStop SQL Installation 1987. - NonStop SQL Programming Reference Manual Reference Manual	Open server designed to provide As many as 255 systems (4,080 users with toolset for building processors maximum) can be custom servers to heterogeneous linked to form a geographically databases. databases.	TANDEM: NonStop SQL
D.S. Patent Number 4,714,995 awarded to TRW on Dec 22, 1987.	Designed to provide the com- plete spectrum of integration needs from automatic updates to query-only ertracts in a heterogeneous data environ- ment.	TBN: TDIR
Purnished upon request.	Continuing work on internal funding, looking for strategic partners to commercialize pro- duct.	XBROX AIT: MULTIBASE

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References	Remarks	
Furnished upon request.	Digital claims Rdb/VMS is the best relational database on the VAX.	DIGITAL BODIPHENT CORP.: Rdb/VMS, DBMS
Purnished upon request.	Claims to excel in personal computer Relational database management, personal computer application development, and host connectivity for access to DB2.	GUPTA TECHNOLOGIES, INC: SQL SYSTEN
Pat Dwyer and Jim Larson "Some Furnished upon request. experiences with a Distributed Database Testbed System", Proc. of IEEE, Vol. 75, No. 5, May 1987. Elmasri et al., "Notes on DDTS", ACM SIGMOD Record, July 1981.	Work redirected to data dic- tionary and engineering tools area due to lack of corporate interest in productizing DDTS.	BOMBYWELL: DDTS
Purnished upon request.	rected to data dic- ind engineering tools the consistency of data values to lack of corporate between related columns of two in productizing DDTS. different tables. DB2 and SQL/ DS participate in Systems Ap- plication Architecture (IBM erclusive) permitting data ac- cess across processor plat- forms. Data from DB2 and SQL/ DS may be accessed/updated from IBM's knowledge based sys- tem products.	IBM: SAA, DB2, AND SQL/DS

References	Remarks	
Purnished upon request.	4GL FOCUS language is the most widely used in the world. IBI claims to be a leader in Hierarchical, Relational, and Network data integration.	INFORMATION BUILDERS, INC.: FOCMET
Purnished upon request.	Designed to handle On-Line- Compler-Applications (OLCP). "Compler" is described in terms of the transaction types and intervals, the datatypes and the transparent network processing features of the database. Ad- vanced as compared with tradi- tional On-Line Transaction- Oriented products.	INTERBASE SOPTWARE CORP.: INTERBASE
Wang, R. and Madnick, S., "Facilitating Connectivity in Composite Information Systems" to appear in ACM Database. Wang, R. and Madnick, S., "The Inter-Database Instance Identification Problem in In- tegrating Autonomous Systems", Proceedings of the Fifth In- ternational Conference on Data Engineering, Feb. 6-10, 1989.	CIS was designed to be non- intrusive and support inte- gration amony very disparate databases. It is part of a larger effort that addresses both advances to theory as well as deeper understanding of the uses and application needs.	MASSACEDSETTS INSTITUTE OP TECHNOLOGY: CIS
Purnished upon request.	Bach database possesses no knowledge of the existence of another database within the Metaphor DIS. The object ori- ented application development environment provides rapid prototype/development capabil- ities to integrate the results of any number of database queries with each other and with data acquired through client developed personal com- puter and VAX based applica- tions.	METAPHOR DIS

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References		Remarks	
	Purnished upon request.	Punded by the Rome Air Devel- opment Center (RADC) and the Maval Ocean Systems Center (MOSC).	BBN SYSTEMS & TECHNOLOGIES: CRONUS
	Purnished upon request.	IDIMS eliminates the need for the user to know query lan- guages, locations of data, database paradigms, or manipu- lation schemata. IDIMS is not tied to any hardware/software platform.	BOO2 ALLEN: Idims
	Purnished upon request.	Surrond SM allows end users to An open protocol will be pub- drive the design process, lished to allow other vendors developers to implement appli- to develop interfaces to data- cations rapidly, eristing base management systems not systems integration to be done on an open systems plat- form.	CAMBRIDGE TECHNOLOGY GROUP: Surround
	Furnished upon request.	An open protocol will be pub- lished to allow other vendors to develop interfaces to data- base management systems not currently planned to be sup- ported by Cincom.	CINCON: Supra

References	Re u arks	
Purnished upon request.	The 4GL provides transaction According to CDC, it is rollback on flat or index file describe the power of AS systems, but makes use of data- base rollback facilities on re- lational systems. An optimis- tic concurrency policy is used to maximize on-line throughput. The 4GL includes a report writ- er, form/screen transaction manager, volume (batch) pro- cessing facility, data diction- ary and other components.	COGNOS CORPORATION: POWERHOUSE
Furnished upon request.	According to CDC, it is hard to describe the power of ASCENT because it is versatile.	CONTROL DATA CORP: ASCENT ¹ suite
The CDM*plus product resulted from Integrated Information Support System (IISS). The first document, in a set of 120, that describes the com- ponents of IISS is: DApro Project IISS, Document Humber: PTR620300001. Obtained at MANTECH DIR, Integration Tech. Div., Attn: D. Judson, WRDC, WPAFB, Ohio 45433.	The CDM*plus enables integra- tion through centralized modeling, rigorous mapping between existing and conceptu- al schema, and robust process- to-process communications capabilities. In addition, it provides a suite of tools for analyzing the impact of change, and bridging to popu- lar modeling tools. It also has extensive reporting facil- base management systems ities.	CONTROL DATA CORP: CDM*PLUS
*plus product resulted M. Templeton, P. Ward, E. Lund, legrated Information "Pragmatics of Access Control system (IISS). The in Mermaid", the IEEE Technical ocument, in a set of Committee on Data Engineering, at describes the com- of IISS is: DApro M. Templeton, "Schema Transla- IISS, Document Number: tion in Mermaid", submitted to 00001. Obtained at conference, Aug. 1989. DIR, Integration Tech. ttn: D. Judson, WRDC, Dhio 45433.	'plus enables integra- rough centralizedDesigned to be modular, port- able, and easy to customize.j, rigorous mapping eristing and conceptu- ea, and robust process- sess communications ities. In addition, ides a suite of tools and bridging to popu- eling tools. It alsoWritten as processes with well- ereplacement of components such as different user interface, output into an application or a database, data dictionaries in ariste reporting facil-blue base management systems.Systems and diverse local data- base management systems.	DATA INTEGRATION, INC.: MERMAID

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Current State Limitations And Puture Directions	Current Installations And Applications	
Puture direction lies mostly in the DB2 and SQL/DS arena, and in complete SAA compati- bility and compliance. bility and compliance.	The RAMIS Information System has approximately 1600 instal- lations in the United States.	ON-LINE SOPTWARE INTERNATIONAL: RAMIS INFORMATION STSTEM
Not all application packages are ported to all 90+ hardware platforms; ongoing effort to resolve this. Specific com- pany directions are to port Oracle across the IBM product offering.	Installations erist within the Department of Defense, Pederal Civilian Agencies, and a wide range of commercial entitities for numerous applications.	ORACLE CORP.: ORACLE
 Puture enhancements include: Automated replication of global model to other nodes. Deferred "snapshot" copies Additional support of unusual native databse data types. Concurrent writes to multi- ple local databases under transaction control (i.e., 2- phase commit). The current INGRES/Star supports multiple- database reads and writes to a single local database under transaction control. 	Sample of customers using INGRES/Star and/or Gateways: Corning Incorporated General Motors Morgan Guaranty 3M Corp. Carnegie Mellon University ICI Americas Scotland Yard Sierra Geophysics	RELATIONAL TECHNOLOGY: INGRES
Software AG provides state of the art tools, implemented us- ing 4th generation technology, which are portable across mul- tiple system environments through its Open Integrated Architecture (ISA).	True distributed processing is currently being implemented at Software AG Darmstadt, Ger- many. Software AG products are currently installed at over 3,000 customer sites. Specifically, in the Pederal Government, Software AG pro- ducts are represented in near- ly 100 agencies and depart- ments.	SOPTWARE AG: Adabas

Current State Limitations And Future Directions	Current Installations And Applications	
Open server is in beta testing and is single-threaded. Con- nection from local database to an open server allows access to all front-end tools available for the Sybase Relational DBMS. Future directions will greatly expand network support, per- formance via multi-threaded processing, and commercial-off- the-shelf gateways to Model 204, DB2, and ANSI-SQL data- bases.	Military Intelligence Organiza- tions. Military Command and Control. Environments similar to the task of integrating TRAMSCOM's subordinate commands.	STBASE: STBASE
Currently implemented only on Tandem NonStop Systems. In- tegration agreements have been signed with Oracle Corp., Sybase, R.T.I. (INGRES), and I.B.I. (Pocus) to have their toolsets ported to Tandem Systems. This will give users access to NonStop SQL through one of these front-ends.	Currently installed in a number The TDI of Portune 500 companies. Ap- plications in various industries of TRM. including finance, manufactur- ing, government, health ser- vices, insurance and marketing. system	TANDEN: NonStop SQL
The TDIE is a TRW commercial product. Currently implement- ing a front-end tool to the TDIE called the Integration Advisor to allow easier construction of the integra- tion interfaces.	E is in production use the Manufacturing Div. It integrates a Planning System on a h a Job Shop Control on an IBM.	TRW: TDIR
Multibase does not currently support updates. Yerox AIT is currently accessing market needs for this capability.	Multibase is installed as a testbed at Cecom at Port Mon- mouth.	XEROX AIT: NULTIBASE

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Current State Limitations And Future Directions	Current Installations And Applications	
Connectivity, front end tool development, and artificial intelligence are ongoing efforts.	POCUS and POCNET are used in financial, academic, and gov- ernment areas, including man- ufacturing, budget tracking, and General Ledgers. There are approximately 4,000 copies of POCUS used on mainframes and micros; approximately 125,000 copies of PC FOCUS have been sold.	INFORMATION BUILDERS, INC.: POCNET
It is not possible to have views which span more than one data- base, or triggers which modify tables in a different database. Work has begun to define a global schema which would allow these features. these features.	InterBase has roughly 1,500 licenses in various industries including aerospace, elec- tronics, financial trading, and manufacturing.	INTERBASE SOPTWARE CORP.: INTERBASE
CIS version 3.0 was completed in August 1989. Certain components have been converted from LISP to C and shell for efficiency. Current CIS supports access to all the databases described earlier. Studies of strategic, organ- izational, physical, and logical connectivity are continuing.	CIS has been used internally at MIT since 8/89 to support the Placement Assistant System (PAS) using the remote data- bases described above.	MASSACHUSETTS INSTITUTE OP TECHNOLOGY: CIS
Metaphor plans to support every popular relational DBMS that offers an SQL Interface. The same or similar iconic access to non-relational databases is under consideration by one or more firms currently providing gateway technology and soft- ware for access to Metaphor's DIS.	been used internally Approximately 10,000 worksta- since 8/89 to support tions in use across 200 govern- cement Assistant System ment and commercial clients in sing the remote data- such as Army Deputy Chief of Labor Statistics, Postal Ser- vice, Transportation Analysis, and Coast Guard for Helicopter Readiness. Also, many Portune 500 companies such as Proctor & Gamble, Chevron Chemical, and British Telcom.	MBTAPBOR DIS

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Current State Limitations And Puture Directions	Current Installations And Applications	
Not currently a commercial pro- duct. System includes documen- tation, tutorials, installation manuals, and telephone hot-line and electronic mail bug report- ing and diagnosis. Puture directions: increased software development toolkit support, additional replication algorithms, more extensive WAN support.	Currently installed at 15 sites Components of IDIMS have been on over 150 machines. Applica- tions include office automa- tion, distributed simulations, expert system decision aides, and an intelligent document retrieval system. C, Portran, and Commonlisp languages are supported.	BBN SYSTEMS & TECHNOLOGIES: CRONDS
See remarks.	Currently installed at 15 sites Components of IDIMS have been on over 150 machines. Applica- developed by Booz Allen for the tions include office automa- tion, distributed simulations, expert system decision aides, and an intelligent document retrieval system. C, Fortran, and Commonlisp languages are supported.	BOOZ ALLEN: Idins
Puture directions include releasing tools to handle compler image as input.	Over 200 different client applications including: Cigna Insurance's Investment Manage- ment Reporting System, Royal Caribbean Cruise Line's CruisePar, and Sheraton's in- hotel reservation system.	CAMBRIDGE TECENOLOGY GROUP: SURROUND
Base software is in general re- lease. Local update supported now. Pirst network implementa- tion in nert maintenance re- lease. lease.	Purnished upon request.	CINCOM: SOPRA

Po an Current State Limitations And Puture Directions Ho de Pl Ho RM	Ov V Current Installations And Applications ag	
PowerHouse StarBase, StarWet, and StarGate were first re- leased in Jan/89 on VAX/VMS and are not yet available on other platforms. VMS-VMS and UNIX- VMS connectivity is available. Interconnection of many other machines is working in the lab, but is not available for re- lease. PowerHouse 4GL cannot be invoked from 3GLs, but can call 3GL subroutines. Power- House StarBase provides embed- ded SQL and GDML support for C, COBOL, and FORTRAN programs (Basic, Pascal and Ada are working in the lab). 3GL ap- plications can access Power- House StarBase, Rdb/VMS and RMS.	Over 12,000 installations on HP VAX and DG employed in a wide range of commercial applica- tions including order proces- sing, inventory control, health care, payroll, financial man- agement, etc.	COGNOS CORPORATION: POWERHOUSE
There has been some thought given to combining the strength of ASCENT with the work done on the IISS system at Wright-Pat- terson AFB. This combination of technologies would replace the current communications methods used by IISS with the *gateway. It would also replace the pre- compiled Cobol programs in the remote hosts with scripts thus removing the requirement for software in the remote host.	ASCENT and LLNL'S TIS software are installed in over 60 places within the government and commercial world. ASCENT was recently included in the AFCAC 251 contract awarded to AT&T by the Air Force recently to be run on the AT&T 3B2- 600 Gs. TSC currently has the TIS version running on several of DLA's DMINS Gould computers.	CONTROL DATA CORP: ASCENT [*] suite
CDM*plus does not support IMS access except through the use of Oracle's IMS/Connect facil- ity which is under develop- ment. The IMS access issue is of major concern within the DOD/Aerospace community. The modeling and the integration of applications and databases is done in an evolutionary manner. Additional DBMS, applications, and computers may be added by user request. Customized implementations are available.	In use at: Northrop Aircraft Division, G.E. Aircraft Engine, US Air Porce Materials Lab., US Navy Air Depot Opera- tion Ctr. Used on the fol- lowing Air Porce Projects: Eng. Change Analysis System, Integrated Sheet Metal Center, Precursor Computer Aided Acquisition and Logistics System.	CONTROL DATA CORP: CDM*PLUS
Mermaid has been recoded to move from prototype to product. Data Integration is now com- pleting support utilities and manuals. Specific DBMS inter- faces may be developed under contract with one user and then licensed to others. Other cus- tomization may be done under contract.	In use at: Morthrop Aircraft Company incorporated 7/89. Division, G.E. Aircraft Mermaid software is owned by Engine, US Air Porce Materials Unisys and licensed to Data Lab., US Navy Air Depot Opera- Integration for product devel- tion Ctr. Used on the fol- lowing Air Porce Projects: Eng. Change Analysis System, Integrated Sheet Metal Center, Precursor Computer Aided Acquisition and Logistics System.	DATA INTEGRATION, INC.: Mernaid

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Approximate Price Range Of Product And Services	Specific Components Of Product Or Service Currently Under Development	Ertent of Changes Required to Local Databases (Hodification Of Data, Programs or Modules Required To Be Resident, etc.)	
tional DBMS products.	The Sybase OPEN-SERVER is in the final stage of beta test- ing. The Sybase multi-level Secure Server is in the final stage of beta testing. This server is the database in DOD's national MLS SYSTEM TESTBED. Support for Oracle is planned, DB2, M204, and IDMSB are under development.	None. (Answer applies to the general case of a heterogeneous data source, whether it is any kind of database management system or not).	STBASE: STBASE
processors: they range from \$4000 to \$64,000 for 1 to 16 processors.	Release 2 of WonStop SQL is cur- rently in test and improves par- allel processing capabilities for performance. Support of nulls will also be provided in that release. that release.	We change is required.	TANDEM: NonStop SQL
and is dependent on the com- plerity of the site specific installation.	- The Integration Advisor is - currently under development with an initial prototype espected in December 1989. Price includes installation	No change is required to any database. An extractor or update capture program may be needed on the source, and an updater program may be needed on the target side.	TRN: TDIB
varying prices.	A reanalysis of the architec- ture for commercialization is under development. Custom software development has	No change is required.	YEROY AIT: NULTIBASE

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Approzimate Price Range Of Product And Services	Specific Components Of Product Or Service Currently Under Development	Ertent of Changes Required to Local Databases (Modification Of Data, Programs or Modules Required To Be Resident, etc.)	
Price ranges depend on the operating environment. A Re- porter ranges for \$19,000 to \$45,000. The full RAMIS In- formation System (including DBMS) ranges from \$49,000 to \$126,000.	All components listed are fully functional and have been released. Continuous enhance- ments are added yearly.	No change is required.	ON-LINE SOPTWARE INTERNATIONAL: RAMIS INFORMATION SYSTEM
Pricing varies greatly based on hardware platform.	Basic software offering is com- plete. Ports to additional hardware platforms is an ongo- ing process. Enhancements to distributed transaction pro- cessing, two phase commit, and global updates for additional platforms.	SQL*Net for network communica- tions; SQL Loader for data import.	OBACLE CORP.: ORACLE
<pre>RMS Gateway \$ 500-\$ 35,000 Rdb Gateway \$1,200-\$ 83,000 INGRES/Star \$1,000-\$ 71,000 INGRES/NET \$ 500-\$ 35,000 Protocol Spt (each) \$ 150-\$ 12,000 DB2 Gateway \$1,500-\$106,000 SQL/DS Gateway \$1,500-\$106,000</pre>	INGRES/Star-global distributed database creation and mainte- nance module. INGRES/Net-communications lay- er supporting communication between heterogeneous hardware via DECNet, TCP/IP, LU 0. LU6.2 support is planned for the near future.	None, ercept to resolve any unusual local DBMS data types that are not supported in INGRES. Por non-INGRES local database management systems, the appropriate DBMS Gateway module must reside on the lo- cal host (eg, the Rdb Gate- way). These Gateway modules are available with INGRES.	RELATIONAL TECHNOLOGY: INGRES
Depending on product, size of machine, and operating system utilized, prices vary from \$700 - \$139,200.	INGRES/Star-global distributedADAMET currently under devel- opment.database creation and mainte- nance module.opment.Planned to be released 11/90.INGRES/Net-communicationslay- via DECNet, TCP/IP, LU 0.VIEW PROCESSOR currently under development.A release date between heterogeneous hardware SOR is only necessary for multiple user views.LU6.2support is planned for the near future.multiple user views.	The database administration or system administration defines all files and views to PREDICT, with unique database identi- fiers. With WETWORK and ADA- MET, no change needs to be made to the applications or the data. To the user, all data access is completely transparent.	SOPTWARE AG: Adabas

Approrimate Price Range Of Product And Services	Specific Components Of Product Or Service Currently Under Development	Brtent of Changes Required to Local Databases (Modification Of Data, Programs or Modules Reguired To Be Resident, etc.)	
Run time only on remote nodes are included in the VMS li- cense. A development license ranges from \$5,000 to \$120,000	Proprietary.	No change is required.	DIGITAL EQUIPMENT CORP.: Rdb/VMS, DBMS
SQLHost - \$20,000. SQLGateway - \$1,995 for each SQLGateway. O. SQLBase single user software de- velopment kit - \$1,295. SQLBase multi user software development kit - \$2,995. SQLRouter - Wo additional charge.	SQLMetwork connections to IBM's SQL/DS on the mainframe and connections to Relational data- base management systems on minicomputers, a UNIX version of SQLBase, a pre-compiler for C, and Lotus 1-2-3 and Excel add-ins for direct access to SQLBase from spreadsheets are under development.	No change is required. SQLBase has local and client-versions that are identical.	GOPTA TECHNOLOGIES, INC: SQL SISTEM
DDTS is still a testhed; not a product. A variety of ser- vices in R&D in this area at available, as long as they are in Honeywell's strategic areas of interest.	Data dictionary tools based on AMSI-standard Information Re- source Dictionary System.	No change is required ercept for distributed updates which require compatible commit protocols.	BONBYWELL: DDTS
Prices depend on model. DB2 - One time charges range from \$113,400 to \$244,850. The monthly lease charges range from \$3,210 to \$4,530. SQL/DS - One time charges range from \$15,980 to \$47,140. The monthly lease charge is \$761. Educational and volume dis- counts are available.	In addition to the pure IBM, non-UNIX solution, IBM is de- veloping a new relational UNIX database. This database will integrate the AIX line (PS/2, RT, S/370), as well as communi- cate with the SAA database line (PS/2-0S/2, AS400, S/370).	Inventory dependent.	IBM: SAA, DB2, AND SQL/DS

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Approsimate Price Range Of Product And Services	Specific Components Of Product Or Service Currently Under Development	Extent of Changes Required to Local Databases (Modification Of Data, Programs or Modules Required To Be Resident, etc.)	
POCNET is approximately \$100,000 for a mainframe, \$50,000 for a minicomputer, and \$1295 for OS/2.	POCNET consists of two types of servers: the data server, and the programmable application server. The data server is available today in all environ- ments. The programmable appli- cation server is available be- tween VAX and IBM, PC and IBM. New implementations will be released in 1990.	No modification to application or data is necessary to enable remote access and cooperative processing. However, if ope- rating system calls are embed- ded in application, these calls moved. moved.	INFORMATION BUILDERS, INC.: FOCHET
Price per license ranges from \$5,000-\$95,000 (0.S. list) de- pending on configuration.	The Oracle gateay is currently in beta test and is scheduled for the release in December 1989 (Oracle on VAX, Apollo, and SUN is planned). All other parts are commercially availa- ble.	The Oracle gateway requires an additional table to be defined in the database, access to some tables in the data dictionary granted, and two servers to be started on the target node. Digital Equipment Corp.'s Rdb requires access to the files through VMS.	INTERBASE SOPTWARE CORP.: INTERBASE
Depends on complexity of the environment.	Many new components are under development, such as: graphical user interface, automated data semantics syn- chronization, semi-automated interdatabase instance identi- fication, LQP generation, interfaces to new databases (LQPs).	None. Local autonomy is a key design criterion. CIS is based on two key design prin- ciples: (1) system non-intrusiveness (2) data non-intrusiveness	MASSACHUSETTS INSTITUTE OF TECHNOLOGY: CIS
\$65,000 - \$ 4 00,000.	DB2, and AT&T 3B2/Oracle gate- way servers (BetaTest). Microchannel workstation com- patibility.	Local autonomy is a key Host Interface language. criterion. CIS is n two key design prin- tem non-intrusiveness a non-intrusiveness	MBTAPHOR DIS

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Approzimate Price Range Of Product And Services	Specific Components Of Product Or Service Currently Onder Development	Extent of Changes Required to Local Databases (Modification Of Data, Programs or Modules Reguired To Be Resident, etc.)	
No license fee for the govern- ment or its contractors. Main- tenance service is approximate- ly \$35,000/year; optional. Ap- plication development support provided at cost.	Extension of network monitoring service, programming support for ADA, security enhancements, performance optimizations, new replication algorithms, porting to new hardware platforms, re- source management algorithms.	No change is required.	BBM SYSTEMS & TECHNOLOGIES: CRONUS
Dependent on application.	Each group of databases and functions must be combined in a different manner, as such Booz Allen has developed systems, such as databases and expert systems, that demonstrate their capability in each area.	No change is required to local databases. IDIMS gueries each of the databases in their native language, based on their struc- ture, semantic and syntactic data stored in the database models knowledge base.	BOOZ ALLEN: IDINS
Por guaranteed fired price delivery of strategic applica- tions prices range from \$150,000 - \$8,000,000 for development. On going mainte- nance is 15% of development cost.	Ertensions to eristing appli- cation network management capabilities. Incorporating ISDN interface to PBX to en- hance eristing telephony capabilities of the environ- ment. Enhance capability to deal with image data.	No change is required.	CAMBRIDGE TECHNOLOGY GROUP: SUBROUND
Licensed on a per CPU basis, varies by CPU size. \$30,000 to \$200.000+.	IMS/DB in detailed design phase. Supra PDM in implemen- tation phase. VSAM completed. Planned local database support: Supra PDM, DB2, RDB, RMS, IMS, DL/1, SQL/DS.	Local DBMS must handle indices. Drop and recreate entity. Local platform must host SUPRA Kernel.	CTNCOM: Sopra

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Approximate Price Range Of Product And Services	Specific Components Of Product Or Service Currently Under Development	Extent of Changes Required to Local Databases (Modification Of Data, Programs or Modules Reguired To Be Resident, etc.)	
First copy is \$5,000 to \$500,000 for Digital VAX. Dis- counts for multiple copies, run-time versions, etc.	StarBase and StarNet on HP3000 and HP9000. Gateways to vari- ous database systems (e.g. Ora- cle). New end-user query tools for decision support. A CASE front-end tool for database de- sign.	Usually none. Some complex data structures may not be readily accessible using Power- House (e.g. graphs, trees, records with variable-length variants, etc.). Once the software is installed, most supported database systems will be readily accessible.	COGNOS CORPORATION: POWERHOUSE
Price is determined by the rel- ative size and power of the UNIX platform running ASCENT. The one-time-charge for a PC- sized platform is \$1,000. The one-time-charge for Super-mini- sized platform is \$30,000.	StarBase and StarNet on HP3000 The DECNet interface currently and HP9000. Gateways to vari- ous database systems (e.g. Ora- cle). New end-user query tools product mode. This capability for decision support. A CASE front-end tool for database de- standard *gateway. sign.	No change is required.	CONTROL DATA CORP: ASCENT*suite
The CDM*plus basic product con- sists of the CDM*plus diction- ary, NDDL, CDM*plus Tools (re- ports, DDL translators, Impact Analysis), NDML, 7 days of training for up to 15 individu- als, and complete installation and support. The cost for the basic product is \$150,000 for a VAX/VMS 8600 version.	Belease 2.0 available in late 1989, provides for full DEC- DEC, and DEC-IBM processing. Peatures available in 1990 include TCP/IP support, addi- tional bridges to CASE tools, erpanded dictionary capability, and additional DBMS access.	The CDM*plus 3 schema architec- None. ture requires only the loading design of the definition and structure of the local database, thus no change is required.	CONTROL DATA CORP: CDM*PLUS
The CDM*plus basic product con- sists of the CDM*plus diction- ary, NDDL, CDM*plus Tools (re- borts, DDL translators, Impact Analysis), NDML, 7 days of training for up to 15 individu- training for up to 15 individu- als, and complete installation and support. The cost for the basic product is \$150,000 for a VAX/VMS 8600 version.	Data dictionary builder and utilities for user installa- tion, recovery, and monitoring are being rewritten. Addition- al DBMS interfaces are being developed. M204 started on government contract, but not completed. DB2 and Teradata interfaces are under develop- ment.	- Kone. Local autono∎y is a key design criterion. e	DATA INTEGRATION, INC.: Mermaid

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Network Services	Global Update And Transaction Management	
Personal computer Networking. LAN support includes: BANYAN VINES, IBM personal computer LAN, Novell NETWARE, and any LAN that is 100% compatible with DOS 3.1. with DOS 3.1.	The RAMIS Data Maintenance Pacility (DMP) provides a way to load files and to modify, browse, and delete data in RAMIS databases and DB2 and SQL/DS tables. A non-proce- dural language is used to ac- complish this. In one request, data may be added, deleted, modified, verified and logged. Transactions may be accepted from various sources such as the screen, sequential transac- tion files, and from the per- sonal computer in fired or comma delimited format.	ON-LINE SOPTWARE INTERNATIONAL: RAMIS INPORMATION SYSTEM
Async, Ethernet, TCP/IP, 3270, STARLAN, LU6.2, DECNEt, CMX, DCAN, NetBIOS, APPC, FTP, EXCELAN, INTERLAN GATEWAY, TOMERNET (Depends upon the hardware platform being used). platform being used).	Offers data and referential integrity for data concurrency, before and after image files and data segment rollback for recovery. Global updates and two-phase commit supported on certain platforms.	ORACLE CORP.: ORACLE
INGRES/STAR via INGRES/NET. Supports thin and thick wire ethernet using DECMet or TCP/ IP protocols, and SMA-LUU.	Current verion of INGRES/STAR supports multi-database reads and updates to a single local database under transaction control. Coming version (beta in Sept., 1989) supports multi-database updates (i.e. 2-phase commit), including full multi-size recovery in the event a local node or the coordinator node should fail.	RELATIONAL TECENOLOGY: INGRES
VTAM, CTCA, DECNEt, Hyperchan- nel, Interlink, TCP/IP and combinations of SMA Gateway.	rrion of INGRES/STAR ADANET supports distributed ulti-database reads transactions which can be a sto a single local series of updates to different inder transaction (beta more than one node within a 1989) supports is received, ADANET synchron- mait), including and checks them for the even- a local node or the tuality of the restart/re- r node should fail. covery procedure.	SOPTNARE AG: Adabas

Network Services	Global Update And Transaction Management	
Supports the standard network protocol for the host environ- ment (e.g. TCP/IP on SUN hard- ware). Committed to erpand network support. Currently provides byte-swapping between protocols supported.	Distributed updates to other Sybase servers available via "two-phase" commit protocol. Does not currently support transactions at the global level; however vendor has com- mitted to do so in the future.	SYBASE: STBASE
A distributed database is im- plemented using an EXPAND net- work which can operate over dedicated lines, leased lines, microwave links, satellite links, X.25 PDNs and Ethernet.	Pull global update and transac- tion management is supported. Any column/row of a distributed database can be updated regard- less of physical location. Transaction Management is pro- vided by the Transaction Moni- toring Pacility (TMP). Each node within the network monitors transactions on the local sys- tem, as the local database transaction completes it sends a message to the originating system which logs a complete transaction after hearing from all nodes.	TANDEM: NonStop SQL
SWA, Ethernet, TCP/IP, MAP/ TOP, Async.	Automatic updates are permit- ted and a cooperative access framework based on "ownership" is utilized. is utilized.	TBW: TDIR
Global Data Manager (GDM) and Local Data Interface (LDI) in- corporate network interface for establishing communication.	Wot supported.	IBROX AIT: NULTIBASE

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Wetwork Services	Global Update And Transaction Management	
DECnet.	Provided by Rdb/VMS and DBMS.	DIGITAL BODIPHENT CORP.: Rdb/VMS, DBMS
NetBIOS, TCP/IP, LU6.2/APPC, SWA/SDLC, all major LANS via NetBIOS, Lan Manager, and LAW Server. Server.	Currently supports multiple databases and multiple cursors per transaction. Two phase commit to be supported in the first quater of 1990 . first quater of 1990 .	GUPTA TECHNOLOGIES, INC: SQL SYSTEM
Communication interfaces based SNA includes: on TCP/IP and MAP protocols - Advanced pe are supported. DSA protocols - Low entry r communications Low entry r Communications Token Ring - Synchronous (SDLC) - X.25.	Transaction management is done using a distributed execution monitor. Updates to certain versions of certain databases are supported.	HOMEYWELL: DDTS
 SNA includes: Advanced peer to peer communications (APPC/LU6.2) Low entry networking (LEN) Token Ring/Local Area Network (TR/LAN) Synchronous Data Link Control (SDLC) X.25. 	Supports: Remote request - Support for single SQL requests. Remote unit of work - Support for several SQL requests within an application program. All requests within a commit scope address a single remote DBMS. Distributed unit of work - Sup- port for several SQL requests within an application program. Each request addresses a single DBMS. The requests within a commit scope can address more than one remote DBMS. Distributed request - Request may address more than one DBMS.	IBM: SAA, DB2, AND SQL/DS

Network Services	Global Update And Transaction Management	
POCNET supports LU 6.2, LU 0, LU 2 (3270 data stream), TCP/ IP, and MetBios. It runs on EtherNet, TokenRing, IBM chan- nels, and VTAM compliant com- munication layers. It is, in the case of LU6.2, compliant with SAA, using APPC and CPI-C facilities. It converts proto- cols when needed; routes mes- sages, controls queuing and task management; authenticates users on networks; works under security subsystems; compresses and encrypts data; and handles error logs.	POCUS utilizes the check/verify model, enabling multiple users to view the same record simul- taneously. The check/verify model tracks the changes of the record by keeping the image of the record in a scratch pad area, and then informs users if the record has been changed.	IMPORMATION BUILDERS, INC.: POCNET
InterBase has a standard message oriented protocol for communica- tions between distributed data- bases. Compatible with TCP/IP, DECnet, and Domain.	Multiple databases can be up- dated in the same interactive session. Prom a program, two additional features, multiple transactions and automatic two- phase commit, can be used to make multiple updates condition- al on the state of individual nodes. Transactions can be tai- lored to commit all, commit some, rollback some, or rollback all data depending on the need of the individual application.	INTERBASE SOPTWARE CORP.: INTERBASE
Uses Ethernet and TCP/IP for Local Area Network access and asynchronous ASCII communica- tions to remote databases. There is a flexible "communi- cations server" module that makes it relatively easy to add new network services.	Present version only supports data retrieval. Puture development will provide various levels of global up- date and transaction manage- ment facilities.	MASSACHUSETTS INSTITUTE OP TECHNOLOGT: CIS
The DIS is LAM-based. It uses Ethernet and XMS for Local Area Network access which can co- erist with TCP/IP, DECnet, and other 802.3 protocols. There are bridges to Token-Ring LAM. WAM communications can be via SNA, APPC internetwork routing, and X.25.	Not supported directly. Meta- phor's object oriented applica- tion development environment can facilitate rapid develop- ment of applications involving small numbers of updates. The DIS includes iconic data entry and modification tools. Direct host based processing is more efficient for high volume transactions.	NETAPEOR SOHOT

Network Services	Global Update And Transaction Manage≣ent	
Synchronous and asynchronous Database queries over a netw representation translation, network monitoring service, LAN are dependent on implementa- and MAN, and dynamic resource location. Underlying transport mechanism is TCP/IP.	Updates my be performed on an object by object basis.	BBN SYSTEMS & TECHNOLOGIES: Cronus
rteer	IDIMS does not update or alter data in any of the databases. Transaction Management is not currently addressed by the sys- tem. tem.	BOO3 ALLEN: IDINS
All network protocols sup- ported including: IBM/SWA (3270, RJE, LU6.2), BSC, TCP/ IP, async, X.25, ISDN. At the application level, Surround is independent of network protocols. protocols.	Global updates performed through Data Handler and com- munication tools on remote systems via database and data- base server in a local data- base. Transaction management function performed by proprie- tary tool, Commander, used for process monitoring, audit trail, and transaction proces- sing.	CAMBRIDGE TECHNOLOGY GROUP: SURROUND
LD6.2, TCP/IP, and DecWet	Secure only when Supra database management systems are in- volved. May be secure for other database management sys- tems that support two phase commit. Insecure for database management systems that do not support two phase commit.	CINCON: SUPRA

Network Services	Global Update And Transaction Management	
PowerHouse StarNet supports DECNet and TCP/IP. 0S/2 LAN Manager has been developed but is not yet commercially availa- ble. MAP support has been prototyped in the lab. prototyped in the lab.	Updates across multiple data- bases are permitted. Power- House StarBase supports the two-phase commit protocol to guarantee the automaticity of such updates. Automaticity cannot be guaranteed for other participants. In practice, most applications can be dis- tributed in such a manner as to avoid the need for consistency across two or more databases.	COGNOS CORPORATION: POWERBOUSE
*gateway supports three types of connections: asynchronous, TCP/IP 802.3 ethernet, and DECMet. The actual network services are provided by UNIX.	This would have to be handled by CDM*plus the DHDMS. *gateway would just and supp be a tool if required to facili- tate the connection to remote ty is en systems. guarant guarant	CONTROL DATA CORP: ASCENT*suite
CDM*plus uses LAW and WAN com- munications within the SNA and DECNet arena on currently sup- ported hardware. TCP/IP is under development and will be available in the first half of 1990.	CDM*plus was designed to allow Transactions that can be passe and support distributed global to one database are permitted. updates. Referential integri- multi-node updates can not be guaranteed. updates can not be ondary. Pragmented relations are updated without guaranteed concurrency. Read concurrency is not guaranteed, but it does make selects and projects of all relations in parallel as a first step to minimize time be tween read of local relations.	CONTROL DATA CORP: CDM*PLUS
s uses LAW and WAN com- uses Uses Ethernet and TCP/IP be- ions within the SNA and tween user interface and Mer- arena on currently sup- hardware. TCP/IP is evelopment and will be to process protocol (LU6.2 le in the first half of planned) between it and the lo- cal database host. All commun- ication between local database translators goes through the controller to avoid many to many protocol problems. Trans- lator may also reside on a UNIX server using TCP/IP to the con- troller, but a special protocol to the DBMS host.	a was designed to allow Transactions that can be passed port distributed global to one database are permitted. Referential integri- Replicated copy updates go to nforced, however, primary copy and then to sec- ode updates can not be ondary. Pragmented relations are updated without guaranteed concurrency. Read concurrency is not guaranteed, but it does make selects and projects of all relations in parallel as a first step to minimize time be- tween read of local relations.	DATA INTEGRATION, INC.: MERMAID

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Query Processing (Retrieval)	Global Schema Construction	
RAMIS Reporter, a comprehen- sive report writer, provides for report development accesi- bility in four ways: non-pro- cedural syntar, menu based, checklist, and ENGLISH (a nat- ural language interface to RAMIS Reporter based on AI principles). All four methods allow for sorting, selecting, manipulating, and formatting of data to produce reports ranging from simple to the more complex production type.	FILE EXPERT, a full screen field description editor, is an integral part of the RAMIS Information System allowing for modification, creation, and deletion of data descrip- tions. Automatic load from dictionaries and COBOL copy books provided where neces- sary.	ON-LIWE SOPTWARE INTERNATIONAL: RAMIS INPORMATION SYSTEM
Once data are imported from the non-Oracle based databases, queries may be performed via the SQL Language. Direct querying by non-Oracle based systems, besides those running DB2 or DBase III, is not per- mitted at this time.	Imported database schemata converted into relational format may be further analyzed to ensure proper naming con- ventions, dependencies and data definitions. No Global Schema exists or is necessary.	: ORACLE CORP.: ORACLE
SQL queries on the Global Dis- All logical files and user tributed Data Dictionary are views are defined in PREDI decomposed into local data- base SQL queries. The statis- fiers. An SQL or WATURAL tically-based distributed query written using this v query optimizer manages the local nodes by developing query execution plans that consider network costs. Thus, through ADAWET tables. interim processing of a com- pler query may occur at the coordinator node. The distri- buted query optimizer enables joins across all the local supported by INGRES gateways.	Global database is constructed ADANET supports replicated manually using SQL to define distributed, and isolated "links" to local databases. Links reference node, local database name, data type, and multiple physical files till support local database table (dataset) as one logical file) throw name. (Future enhancements include automated maintenance of global schema to reflect changes to local schemata, and to other nodes).	RELATIONAL TECHNOLOGY: INGRES
ies on the Global Dis- Data Dictionary are ed into local data- queries. The statis- fiers. An SQL or WATURAL based distributed timizer manages the timizer manages the the appropriate database identi- query written using this view vill be dynamically routed to the appropriate database at ecution plans that network costs. Thus, processing of a com- ry may occur at the tor node. The distri- ery optimizer enables d by INGRES gateways.	Global database is constructed ADAWET supports replicated, manually using SQL to define distributed, and isolated "links" to local databases. Links reference node, local multi-file user views (ie database name, data type, and multiple physical files viewed local database table (dataset) as one logical file) through a name. (Future enhancements product called VIEW PROCESSOR include automated maintenance currently under development. of global schema to reflect changes to local schemata, and to other nodes).	SOPTWARE AG: Adabas

Query Processing (Retrieval)	Global Schema Construction	
A query in SQL is passed di- rectly to the SQL server. All other queries are executed as parameters to a stored proce- dure which triggers an event handler to a request written in the native database lan- guage which accesses the local database.	On-line active data dictionary maintains relational model for Global Schemata. Local data- base schemata, subject to schema and data translation mechanisms, are maintained in Global Schema in relational form.	STBASE: STBASE
SQL queries are processed on the local database system. If a query consists of data from a distributed database, each sys- tem processes the query against the local database and transmits only the selected columns/rows to the querying system for final processing.	The idea of a Global Schema does not exist. Users create local schemata and can access remote schemata with standard data manipulation language. Users have access to distributed data- bases and can execute distri- buted joins to access a logical- ly local database.	TANDEM: NonStop SQL
A local database SQL query can be processed against a remote IMS database (retrieval only).	A front Integrat interact global j	TRM: TDIR
database SQL query can A query expressed in DAPLEX on ssed against a remote Global Schema is decomposed in- base (retrieval only). to queries on DAPLEX Local Schemata. These are further translated into Local DML quer- ies. Adopts "no two variables in a query range over the same entity" strategy to generate queries on Local Schemata. Adopts data reduction and data movement techniques in Global Query optimization strategy. Merging and final processing takes place at Global Schema result node.	end tool called the Uses the Functional Data Model cion Advisor allows the to define the global schema and ive construction of a DAPLEX as the GDML. The local integrating schema. DBMS and their local host sche- ma which are expressed in terms of the Global Data Model. The local schema are merged into global schema are merged into ntegration Schema with an auxiliary Integration Schema which con- tains information for reconcil- ing inconsistencies between local schemata.	XEROX AIT: MULTIBASE

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Query Processing (Retrieval)	Global Schema Construction	
Query processing is done through an optimizer which takes into account Table Car- dinality, Index Cardinality and type of Query. Directions on how to process gueries can be specified in Network Manip- ulation Language (NML).	Uses CDD/Plus dictionary.	DIGITAL BODIPMENT CORP.: Rdb/VMS, DBMS
Pull SQL support for queries from various query tools (SQL- Talk, WindowsTalk, Express Win- dows, APIs, SQLWindows, etc.). Read only capabilities exists ir database. ities exists ir	Supports full Relational cata- log. Several tools for defini- tion of database design are available. available.	GDPTA TECHNOLOGIES, INC: SQL SYSTEM
E-C-R based GORDAS is decom- posed into internal, rela- tional algebra based guery trees. These query trees are then optimized based on local processing and network costs into an execution strategy with the right communication commands embedded in the query tree. A SQL guery interface is supported. Also, SQL queries will be converted into GORDAS gueries first. The sub query-trees are converted at each local site into local DBMS commands.	Global schema is constructed A single global schem manually and populated using exist, yet many compo dictionary files into the data bases may be defined. structures of the system.	HOMBYWELL: DDTS
Queries are optimized for mini- mal communication cost. Quer- ies can be submitted against any number of composite data- bases provided the user is authorized to do so.	A single global schema does not exist, yet many composite data- bases may be defined.	TBM: SAA, DB2, AND SQL/DS

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	Global Schema Construction	Query Processing (Retrieval)
INPORMATION BUILDERS, INC.: Pocnet	Several facilities exist in the Global Schema does not exist. POCUS language to create schemata. One is File-Talka window driven facility to in- teractively build an MFD. Other examples are AUTOSQL and AUTOIDMS which directly take data information from data- base catalogues. The ACE (Application Control Environ- ment) reconciles application and security requirements across the network.	The POCUS language can directly A query expressed in SQL is de- access the aforementioned data- bases. It can also dynamically against local schemata. A pro- create views of new fields at the database level or summation if the database is local or re- level. Processing is optimized mote and routes it through the based on the functionality of the database, which includes record-by-record and set pro- cessing. eessing. database of Oracle, it is restored to SQL through the gateway and passed to the Oracle access method.
INTERBASE SOFTWARE CORP.: INTERBASE	Global Schema does not erist.	US language can directly A query expressed in SQL is de- the aforementioned data- composed into BLR queries It can also dynamically against local schemata. A pro- views of new fields at gram called Y-valve determines abase level or summation if the database is local or re- Processing is optimized mote and routes it through the functionality of network if remote. In the case abase, which includes of InterBase or Digital Equip- by-record and set pro-
MASSACEUSETTS INSTITUTE OF TECHNOLOGY: CIS	ingle global sche- ple user-oriented schemas. The a is created by administration lited and modified	Query is parsed and validated against the global schema. Optimal sequence of steps is determined to restrict local queries, whenever possible. Ability for user to provide rules to override normal optimization on even inter- actively influence sequence. Each LQP handles requests for data from its database, as decomposed from the user's request by the Global Query Processor (GQP).
METAPHOR DIS	a single global sche- local schemata of one or more databases are graphically de- ion schemas. The chema is created by base administration be edited and modified join) across two or more physical databases, Metaphor's object oriented development environment (the capsule) can be used.	A query is performed by point- and clicking at objects pre- sented as a logical view tai- lored for the specific end user. End user menus, browse lists, and computed fields are presented in a similar manner. SQL queries are automatically generated by the DIS tool set. Syntactical differences between the various flavors of SQL are handled transparently by the gateway server.

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Query Processing (Retrieval)	Global Sche∎a Construction	
Query processing (QP) across distributed object-oriented data storage is supported; QP over object-oriented repli- cated data is under develop- ment. ment.	Global Schema definition is supported for object-oriented data storage and retrieval. Schema defined when specifying the interface definition for an object type.	BBN STSTENS & TECHNOLOGIES: Cronus
High level query is broken into a series of low level queries and function calls by the query processor (a constraint satis- faction knowledge based system to generate a plan to resolve the high level query. Knowledge from the Database Knowledge base and the Functions Knowledge base is used to retrieve and alter data). Mert, each query is called by the Query Language Generator (QLG). The QLG may be spawned to query databases si- multaneously to enhance effi- ciency.	Not maintained beyond those in the Database Models Knowledge Base.	BOOS ALLEN: Idims
Query generator receives re- quest for data, generates guery, and routes to appro- priate data servers and remote systems. systems.	Data Handler (patent pending) canonicalizes all data into a form that facilitates inde- pendence of user interface, functionality and data sources.	CAMBRIDGE TECENOLOGY GROUP: SURROUND
A global task running in the global layer spawns local tasks running in the local layers of each Supra involved in the query. Local tasks are subdi- vided to local database manage- ment systems by the HDMP. ment systems by the HDMP.	Defines entities to Global Directory. These are validated at definition time against the local schema of the target DBMS. BBMS.	CINCON: SOPRA

Query Processing (Retrieval)	Global Schema Construction	
Joins that involve tables from different databases are resolv- ed by simple nested-loop re- trieval. Access to relational systems is performed by gen- erating suitable DML and using the most efficient interface mechanism available (e.g. Rdb/ VMS DSRI, DG/SQL HLI). Re- search on a distributed query engine is in progress.	There is no global schema per se. Multiple relational schemata may be referenced within the same program. Table names must be qualified in DML statements by a logical data- base name.	COGNOS CORPORATION: POMERHODSE
• Queries would be written in a - combination of C and Shell Scripts. The query would call + gateway and the script would extract the specific data re- tion. /	The DHDBMS would handle this area.	CONTROL DATA CORP: ASCENT*suite
CDM*plus relational operations (query, update, insert and de- lete) are based on relational views as created in the exter- nal schema described above. The query is embedded in cur- rent applications replacing multiple queries and proces- sing or it is created as part of new application develop- ment. All DBMS specific code is generated regardless of location and DBMS type.	CDM*plus dictionary/directory uses the ANSI/SPARC 3 schema architecture to provide global schema mappings. This allows full access to all enterprise data descriptions and ac- comodates mapppings between schema.	CONTROL DATA CORP: CDM*PLUS
CDM*plus relational operations Query in SQL is parsed and val- (query, update, insert and de- lete) are based on relational An optimizer selects local views as created in the exter- anal schema described above. The query is embedded in cur- rent applications replacing multiple queries and proces- sing or it is created as part of new application develop- ment. All DBMS specific code is generated regardless of location and DBMS type. In the translate for one database. In the translator does schema, data, and language translation.	Many global schemata may exist for different applications de- lfined across various local databases. Definition done with an off-line tool. Local schemata are loaded. Each field is labelled with a domain (semantic labelling). Database administrator is assisted in specifying replicated and fragmented relations and in resolving different representa- tions for data elements of the same domain.	DATA INTEGRATION, INC.: MBRMAID

Data Incompatibilities And Semantic Mismatches	Local Database Schema Conversion	Global Data Model	
Performs format, length and data type conversions where necessary. Allows for usage and actual data format types. Conversion occurs during usage, so there is no need to change eristing data.	Conversion is not required. RAMIS interfaces are designed to support databases 'as is'. Definitions are automatically loaded from the dictionary or catalog (e.g. DB2 catalog or IDMS - IDD) to the RAMIS data dictionary. A full screen editor provides modification functionality.	Relational.	OR-LINE SOFTWARE INTERMATIONAL: RAMIS INFORMATION SYSTEM
Supports different data types and units assigned to seman- tically equivalent fields through data definition in the Data Dictionary. Data conversion is performed as the data are physically imported.	Non-Oracle Based schema is converted into a relational model via software utilities and SQL*Loader software which remaps imported data into Oracle format. Base relations and semantic information map- ping is done before data are actually brought into Oracle.	Relational.	ORACLE CORP.: ORACLE
Performs data type and error message conversions to a gene- ric set, data dictionary conversions via a standard dictionary views interface; resolves syntactic and seman- tic mismatches of SQL ("Open SQL") for relational local databases, and via translation of SQL to local database lan- guage for non-relational DBMS.	None. Global schema is a set of links to local schemata. The link points to a local database table, or to individ- ual fields (columns) in a local database table. Links observe the protections and integrity controls that local database administrators estab- lish for their data.	Relational.	RELATIONAL TECHNOLOGY: INGRES
There is none. ADANET acts as a translator.	obal schema is a set ADAMET defines and processes to local schemata. one or more physical databases points to a local at a logical level. ADAMET table, or to individ- maintains a mapping of logical identifiers or user views, abase table. Links with the corresponding physi- he protections and cal database/file identifiers controls that local which comprise the logical administrators estab- their data. parency is provided.	Relational-like.	SO ptuare ag: Adabas

Data Incompatibilities And Semantic Mismatches	Local Database Sche∎a Conversion	Global Data Model	
Local database schemata conver- ted to the relational form are placed directly in Global Schema. Functional transla- tions must be programmed in the open server. Format conver- sions are handled automatical- ly. Supports user defined data types for global manipulation language.	Weeds conversion into rela- tional form or information may be passed as the result of a preprogrammed event written in the local database language.	Relational (Entity-Relationship Relational based).	STBASE: STBASE
Data type and format are defined Maintains information in a in local schemata which are knowledge base that allows accessible from remote systems. level of semantic check to applied through the use of rules or programs.	No conversion is required. Non- Transfor Stop SQL is designed to support into glo a distributed database on Tandem one with NonStop Systems. integrat	Relational.	TANDEM: NonStop SQL
d Maintains information in a knowledge base that allows any level of semantic check to be applied through the use of rules or programs.	 Transformation can either be into global model or one-to- one with the aid of a global integrating schema. 	BCR. Also allo⊌s ∎ultiple global schemata.	TRN: Tdir
<pre>s information in a Resolved through an Integration in a patabase Schema (called IS) for all the local schemata. Extent through the use of of coverage is not sufficient, but it can take care of differences like kilometers and miles, age in years and in qualitative form, (young, old), etc.</pre>	Local database schema called Local Host Schema (LHS) in any model needs to be redefined completely in the Punctional Data Model to enforce uni- formity. New schema is called Local Schema (LS). Local host schema remains intact.	Functional Data Model.	IEROI AIT: Multibase

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Data Incompatibilities And Semantic Mismatches	Local Database Schema Conversion	Global Data Model	
Performs data conversion when possible.	The schema provided by DBMS is sufficient for conversions.	Wetwork and Relational	DIGITAL BODIPMENT CORP.: Rdb/VMS, DBMS
Automatic Import/Export is pro- vided for: Lotus 1-2-3, dBase 2/3, CONDOR, Enable, DIF, and ASCII files. SQLBase has auto- matic data conversion for "Un- load/Load" to DB2.	No conversion is required. Local database schema is the same as client-server schema and IBM's DB2 schema.	Relational (ertended)	GUPTA TECENOLOGIES, INC: SQL SYSTEN
Mappings between local data- base schema and the canonical schema are stored in the data dictionaries.	Local database schemata need to be converted into E-C-R model and integrated with each other manually. Using these converted and integrated sche- mata, the global and local data dictionaries will be pop- ulated.	E-C-R and Relational	HOMBYWELL: DDTS
Supports a wide range of data types: integer, decimal, floating point, character, graphic, and date/time. A Data Extract Facility allows data to be extracted from other databases (including non-IBM: ADABAS, M204, IDMS, DATACOM/DB, and NOMAD2), and files for use by DB2 and SQL/DS.	Conversion to relational form is required.	Relational	IBM: SAA, DB2, AND SQL/DS

Data Incompatibilities And Semantic Mismatches	Local Database Schema Conversion	Global Data Model	
Each database has corresponding interfaces, even between SQL databases, such as DB2, Oracle, Ingres, and Sybase, Through the interface and the MPD, all data and semantic incompat- ibilities are reconciled.	Each local database has corre- sponding schema, called a Mas- ter File Description (MPD). FOCNET interfaces with DSRI and other systems directly. In turn, some databases, like Dbase III or Rdb, do not re- quire a FOCUS Master File De- scription. In all cases, a Master File Description can be created automatically.	Hierarchical, Relational, and Network.	INPORMATION BUILDERS, INC.: POCNET
Data format incompatibility be- tween heterogenous systems is transparently handled by an XDR layer.	Weeds conversion to relational form. Maps into existing data dictionary and extends it when necessary.	Relational.	INTERBASE SOPTWARE CORP.: INTERBASE
Major emphasis on supporting joins between disparate data- bases through synonym and translation facilities. Tools to aid in inter-database instance identification (through key semantic match- ing, and organizational af- finity) are being developed.	Entity and attribute restruc- turing: many-to-one and one- to-many mapping (both local to global and global to local). Automatic invocation of joins and projections, as needed.	Ertended ER.	MASSACHUSETTS INSTITUTE OF TECHNOLOGY: CIS
Data from any respective data- base are translated into a com- mon form within the Metaphor DIS environment. Syntactical differences between the various flavors of SQL are handled transparently by the gateway server.	and attribute restruc- many-to-one and one- tionary (WTDD), stored on each mapping (both local to database/host computer, is used and global to local). by the graphical Query, Data- ic invocation of joins base Administration (DBA) and jections, as needed. Reporter Tools to organize the physical database tables into logical views for end-user, developer access.	Relational.	MBTAPEOR DIS

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	Acrony m Derived	Environment (Software and Hardwa	Local Data Modela Sup	Local DBMS Supported	Global D Manipula Language
	Acronym Derived Prom	nvironment Software and Hardware)	Local Data Models Supported	DBMS rted	Global Data Manipulation Language
ON-LINE SOPTWARE INTERNATIONAL: RAMIS INFORMATION SYSTEM	Rapid Access Management Infor- mation System	General. RAMIS runs in the following operating environments: DOS/ VSE, DOS/CICS, VM/CMS, MVS/TSO, MVS-XA/TSO, MVS/BATCH, MVS/ CICS, MVS-XA/CICS.	Hierarchical, Relational, and Network.	DB2, IDMS/R, ADABAS, SQL/DS, TOTAL, IMS/DB, TERADATA, DATACOM/DB, DL-1, VSAM, ISAM, and QSAM.	RAMIS non-procedural query language. An English-like syntar is also available through a complete menu-based architecture.
ORACLE CORP.: ORACLE	None	General. Oracle relational DBMS is cur- rently supported on over 90 hardware platforms.	Relational.	DB2, SQL/DS, and DBase III.	SÕT
RELATIONAL TECHNOLOGY: INGRES	None	General. VMS and Unir platforms.	Relational, Hierarchical, and Network.	Version 5: Rdb and VAX Record DL/1, VSAM, TOTA Management Service file system. SQL/DS, and DB2. Version 6: DB2, Rdb, RMS, IMS, and SQL/DS.	SQL. Also a forms interface for defining global model.
SOPTHARE AG: Adabas	Adaptable Database Management System	General. Production, development, and information center environ- ments. Software AC's product line supports multiple hardware configurations and operating systems.	An inverted list structure with relational characteristics. Data are stored completely sep- arate from the relationships. Both hierarchical and network structures are supported.	DL/1, VSAM, TOTAL, INS/DB, SQL/DS, and DB2.	ADABAS Direct Calls, ADABAS SQL (SQL embedded in COBOL, PORTRAN, or PLI), or NATURAL (4th Generation Language).

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	STBASE: STBASE	TANDEM: NonStop SQL	TRM: TDIE	XEROX AIT: Multibase
	System Database	None	TRW Data Integration Engine	None
Acronym Derived Prom				
Rovi rooment	General. General. VAX. STN. Apollo. Pyramid.	General. Designed to run on Tander Won-	General. General. The TDIR currently rung on any	General. Multibase runs on any operating
(Software	Stratus, IBM and various other	Stop systems.		system supporting ADA.
and Hardware)	platforms supported.		UNIX/VMS.	
Local Data Models Supported	Hierarchical, Relational, and Retwork via Relational Inter- face.	Relational.	Hierarchical, Relational, and Network.	Hierarchical, Relational, Net- work, and Pile Systems.
Local DBMS Supported	Sybase.	NonStop SQL and Tandem NonStop Systems.	IDMS, IMS, DB2 on IBM. Oracle on VAX. Sybase on SUN.	IDS II, S2000, Oracle, and RIM
Global Data Manipulation Language	Transact-SQL, an extension to AMSI SQL, permits transaction management.	ANSI SQL plus ertensions.	Wone required.	DAPLEX. The local host may not support all capabilities pro- vided by DAPLEX.

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Data Incompatibilities And Semantic Mismatches	Local Database Schema Conversion	Global Data Model	
Data representation conversions between different architectures are supported.	Mone required.	Object oriented.	BBN SYSTEMS & TECHNOLOGIES: Cronus
Resolved through the use of functions modelled in the Punc- tions Models Knowledge Base. The frames in the knowledge base have a representation of a func- tion and how data are altered or manipulated by that function. Punction applications, as well as database gueries, are deter- mined by the Query Processor's constraint satisfaction proce- dure.	Local database schemata are con- None required verted into frame network in Database Models Knowledge Base.	Prame Based/Object Oriented. Supported through the Database Models Knowledge Base. The Query Language Generator uses information derived from the Knowledge Base to determine the appropriate query language and machine.	BOOZ ALLEN: Idins
Total system and data integri- ty guaranteed via canoni- calized Data Handler (patent pending).	- None required.	Por each remote system access is made either via application on that system, database query, or direct file access. In short, any technique that does not require code to be written for remote system.	CAMBRIDGE TECENOLOGY GROUP: SURROUND
Handles semantics with exten- sions to the CREATE statement as well as supports popular data formats.	Generic CREATE statement within the Global Data Mode. SUPRA offers extensions to the CREATE statement to handle semantic and conversion issues.	Relational	CINCOM: Supra

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	COGNOS CORPORATION: POWERHOUSE	CONTROL DATA CORP: ASCENT'suite	CONTROL DATA CORP: CDM*PLUS	DATA INTEGRATION, INC.: Mermaid
Global Data Model	Relational with ertensions for arrays, substructures, redefin- itions and network relation- ships.	See above.	A 3-schema architecture with conceptual schema stored in IDEPir, and ER based modeling method.	Relational.
Local Database Sche≣a Conversion	No local conversion supported. A proprietary schema repository (dictionary) is provided for systems which do not have schema data (such as flat or indered files). Relational schema data are taken directly from the relational DBMS system tables and mapped into canoni- cal form by database specific routines.	The programs written to commun- icate with the remote systems would have to know what the format of the data being er- tracted was and what the target format needed to be before it can make the conversion.	Local database schemata remain intact and are described to CDM*plus as Relational, Net- Work or Hierarchical schema.	Structure conversion: One to Many relation mapping (global to local) with joins added, Many to One relation mapping using different projections, and One to Many and Many to One field mapping. Model conversion for network model: Relation (global) to repeating group and Joins (global) to links.
Data Incompatibilities And Semantic Mismatches	The data model describes most sequential, relative, and in- dered files, as well as most relational and network data- bases. A large number of data types are supported providing 1:1 coverage. Conversions be- tween different data types of the same kind are performed automatically.	See above.	Through the 3 schema defini- tions and inter-schema map- pings, CDM*plus performs all the necessary data format and unit conversions across the schemata. CDM*plus NDML in- cludes full referential in- tegrity and domain constraint checking.	Deals with two types of data translation: functional type and enumerated type. Func- tions, written in C, deal with problems like unit and format conversions. Enumerated type translations deal with convert- ing sets of values through an encode/decode table lookup. Translation is done on values in query and on data when a re- lation is first accessed.

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