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#### **Software Managed Cache for Parallel Systems**

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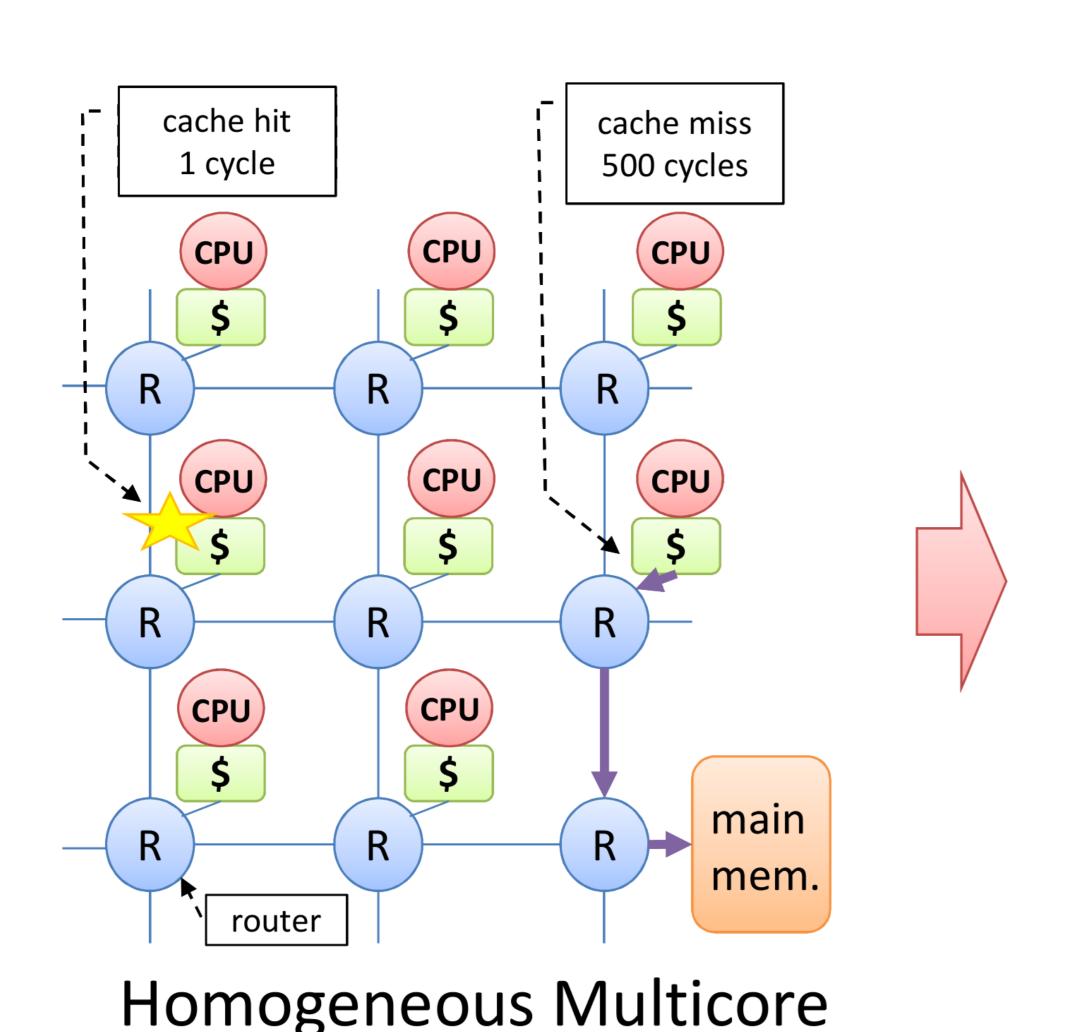
# Software Managed Cache for Parallel Systems

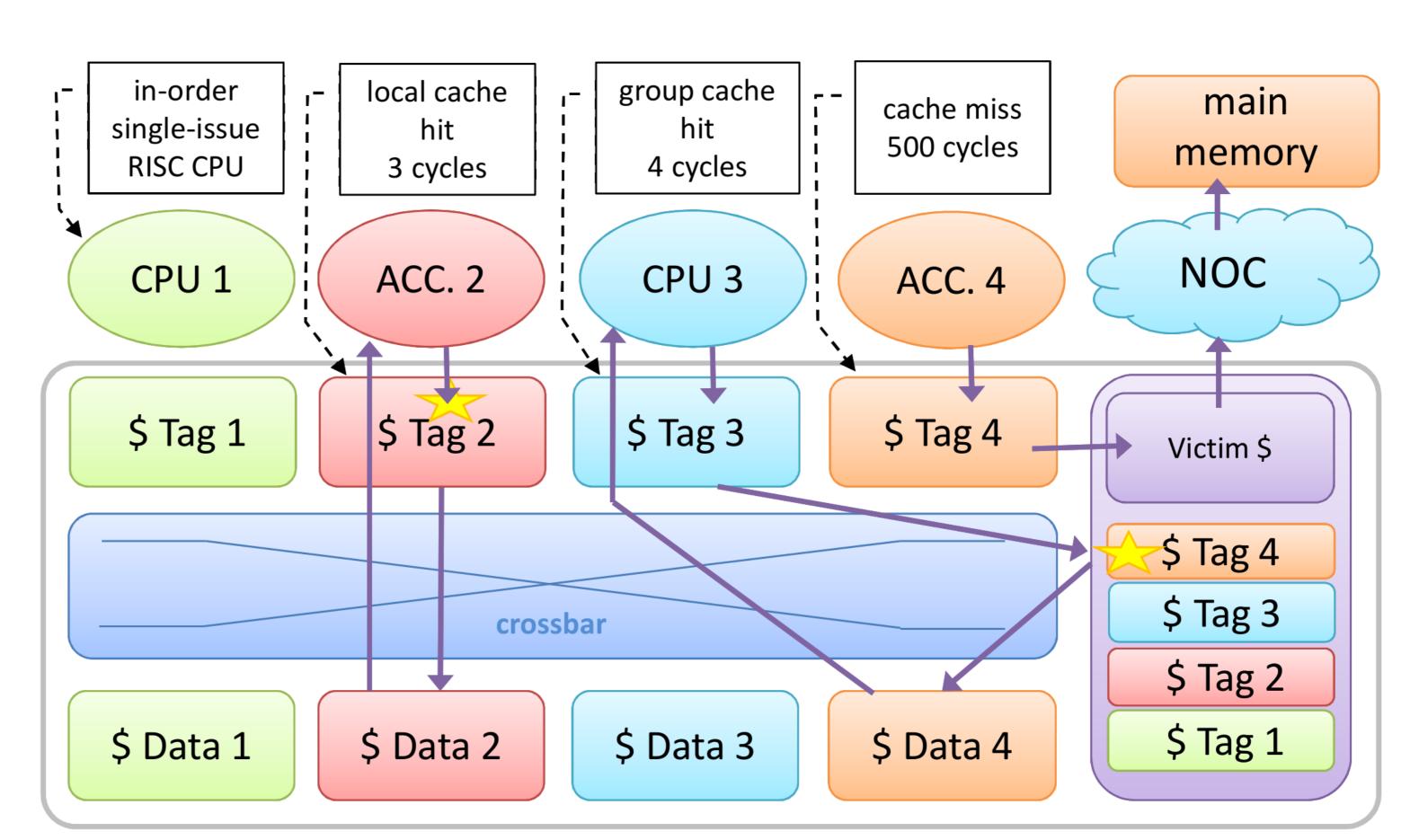
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# Current Systems

# **Envisioned Future Systems**





Heterogeneous Multicore

### Motivation

- ► Homogeneous multicores tend to have small private caches
- ► We envision heterogeneous multicores with a shared cache
  - > can benefit from data stored in neighboring cores
  - ► cheap memory consistency within a group of processing elements
  - ▶ data traffic on the global interconnect is reduced
- ► We argue for a software managed cache!

### Contributions

- ► Software managed multi-banked first level data cache
- ► Application aware software controlled replacement strategies

# Implementation

- ▶ Use a hardware efficient and energy efficient 4-way set associative cache
- ► Cache hit check is done in the following way
  - ▶ tags are checked in sequence, one tag per cycle
  - ▶ a hashing function is used to predict what tag to check first
  - we end when we find a hit or there are no more tags
  - ▶ this leads to increased associativity at a low power consumption
  - ► however, unless the first tag checked is a hit, the cache hit time is increased
- ► Use hardware and software to implement replacement policy
  - on cache misses, CPUs can execute a cache replacement algorithm
  - ▶ balancing cache by relocating cache lines
  - replacement policy may change dynamically
  - ► try to avoid the software replacement algorithm when the expected memory latency is low
  - ▶ use a simple algorithm implemented in hardware in cases where the software replacement algorithm is too slow
- ► Specific memory regions can be labeled
  - often used variables may have a higher priority
  - > specific memory regions may have preferred locations in the cache
  - > specific memory regions may be locked in cache
- ► Additional costs: duplication of cache tags

# Replacement Policy

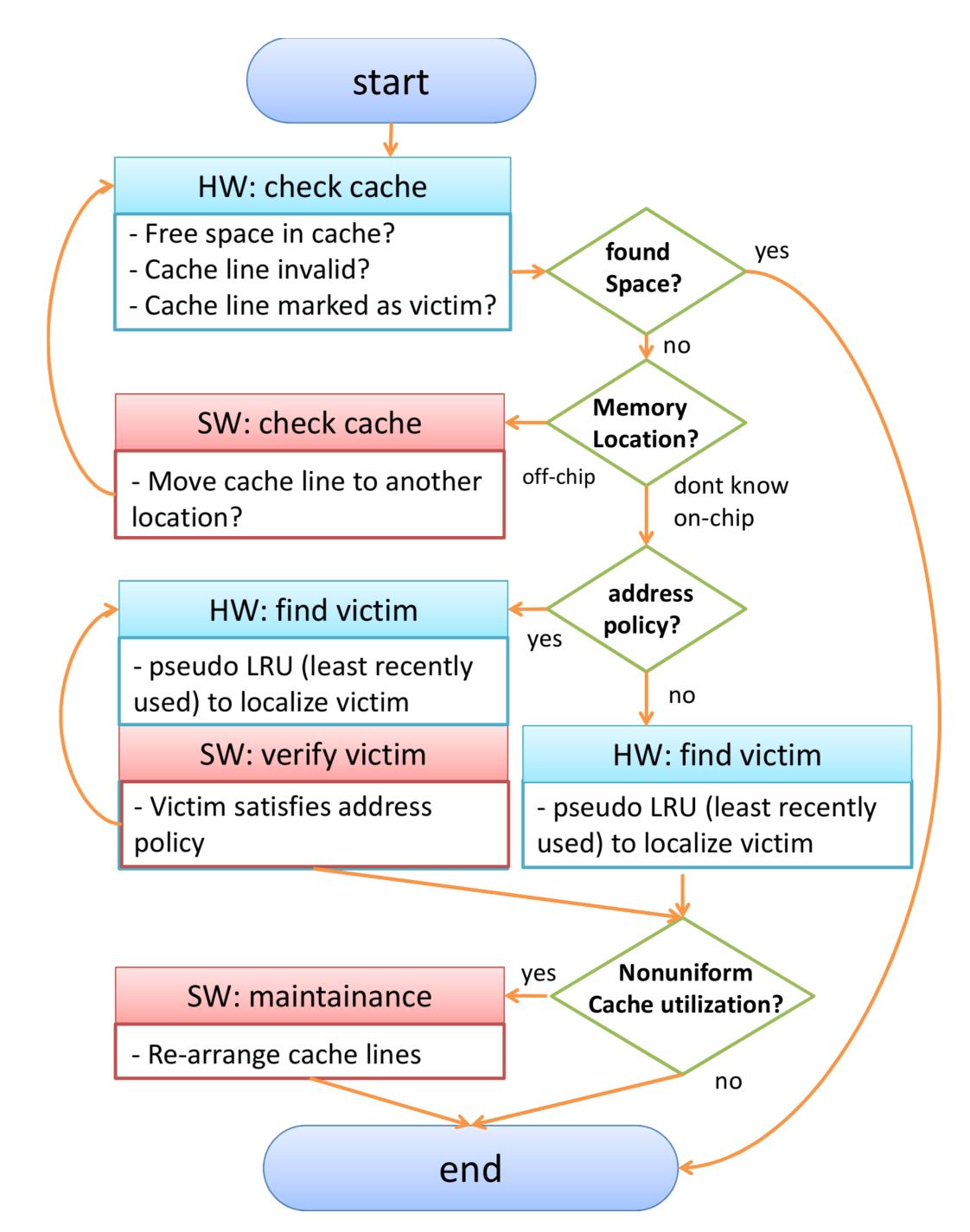


Figure: replacement policy flow graph

Depending on where memory is fetched up to hundreds of clock cycles are used. Embedded processors typically stall on a cache miss.

Instead of waiting for the cache miss to be resolved processors can execute an advanced cache replacement algorithm.

### Conclusions

- ► We propose a software managed multi-banked first level data cache for parallel systems
  - highly configurable
  - ► more area and power efficient than a pure hardware implementation of highly associative cache
- ► We propose an application aware software controlled replacement strategy
  - ▶ use both hardware and software to implement replacement policy