

MODIFIED CHEBYSHEV-PICARD ITERATION METHODS FOR SOLUTION
OF INITIAL VALUE AND BOUNDARY VALUE PROBLEMS

A Dissertation

by

XIAOLI BAI

Submitted to the Office of Graduate Studies of
Texas A&M University
in partial fulfillment of the requirements for the degree of

DOCTOR OF PHILOSOPHY

August 2010

Major Subject: Aerospace Engineering

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ABSTRACT

Modified Chebyshev-Picard Iteration Methods for Solution of Initial Value and
Boundary Value Problems. (August 2010)

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The solution of initial value problems (IVPs) provides the evolution of dynamic system state history for given initial conditions. Solving boundary value problems (BVPs) requires finding the system behavior where elements of the states are defined at different times. This dissertation presents a unified framework that applies modified Chebyshev-Picard iteration (MCPI) methods for solving both IVPs and BVPs.

Existing methods for solving IVPs and BVPs have not been very successful in exploiting parallel computation architectures. One important reason is that most of the integration methods implemented on parallel machines are only modified versions of forward integration approaches, which are typically poorly suited for parallel computation.

The proposed MCPI methods are inherently parallel algorithms. Using Chebyshev polynomials, it is straightforward to distribute the computation of force functions and polynomial coefficients to different processors. Combining Chebyshev polynomials with Picard iteration, MCPI methods iteratively refine estimates of the solutions until the iteration converges. The developed vector-matrix form makes MCPI methods computationally efficient.

The power of MCPI methods for solving IVPs is illustrated through a small perturbation from the sinusoid motion problem and satellite motion propagation problems. Compared with a Runge-Kutta 4-5 forward integration method implemented in

MATLAB, MCPI methods generate solutions with better accuracy as well as orders of magnitude speedups, prior to parallel implementation. Modifying the algorithm to do double integration for second order systems, and using orthogonal polynomials to approximate position states lead to additional speedups. Finally, introducing perturbation motions relative to a reference motion results in further speedups.

The advantages of using MCPI methods to solve BVPs are demonstrated by addressing the classical Lambert's problem and an optimal trajectory design problem. MCPI methods generate solutions that satisfy both dynamic equation constraints and boundary conditions with high accuracy. Although the convergence of MCPI methods in solving BVPs is not guaranteed, using the proposed nonlinear transformations, linearization approach, or correction control methods enlarge the convergence domain.

Parallel realization of MCPI methods is implemented using a graphics card that provides a parallel computation architecture. The benefit from the parallel implementation is demonstrated using several example problems. Larger speedups are achieved when either force functions become more complicated or higher order polynomials are used to approximate the solutions.

Dedicated to my parents
for their unconditional love and always being there

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CHAPTER I

INTRODUCTION

Improved methods for solving initial value problems (IVPs) and boundary value problems (BVPs) are of fundamental importance for analyzing and controlling dynamic systems that are described by ordinary differential equations. IVP solutions generate trajectories showing how the dynamic system evolves with time with the given initial conditions; whereas BVP solutions give the system trajectories that satisfy several constraints defined at different times. Although there already exists a large literature for solving IVPs and BVPs, there remain several challenges and therefore a need to improve the current methods and develop new approaches. Among the issues are: (i) How to significantly reduce the computational burdens; and (ii) How to optimize the approaches to utilize emerging parallel computing architectures.

This dissertation proposes a modified Chebyshev-Picard Iteration (MCPI) methodology to solve IVPs and BVPs. Compared with most existing methods (that are based on the forward numerical integration methods) to solve IVPs and either the direct or indirect methods (see definitions in Section 3) to solve BVPs, the proposed MCPI methods are designed for computational parallelization so they can take advantage of the rapidly developing parallel technology. Applications of MCPI methods to several celestial mechanics problems are explored. The studies show that the proposed MCPI methods can achieve both computational efficiency and high accuracy, even prior to parallel implementation. One challenge of the proposed MCPI methods, the limited convergence domain, is addressed by proposing several techniques that are shown to enlarge the convergence domain. Using a Graphics Processing Unit (GPU)

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architecture, we discuss one parallel implementation of MCPI methods and illustrate the benefit of using this approach.

The remainder of this chapter is organized as the follows: In Section A, existing parallel algorithms used to solve IVPs and BVPs are discussed, followed by a discussion of historical literature on using Picard iterations, Chebyshev polynomials, and Chebyshev-Picard methods for solving IVPs and BVPs. In the context of the existing Chebyshev-Picard methods, the unique features of the proposed MCPI methods are then summarized. Section B presents several IVP and BVP application fields to which the methods developed in this dissertation are applicable.

The remainder of this dissertation is arranged as follows: Chapter II introduces formal statements of IVPs and BVPs that are the focus of this dissertation and develops MCPI methods that can be modified to solve both IVPs and BVPs. A vector-matrix form of MCPI methods is then presented. The related fundamentals of the classical Picard iteration and Chebyshev polynomials are summarized in Appendix A and Appendix B, for the readers' convenient reference. Chapter III focuses on applying the MCPI methodology to solve IVPs. A piecewise approach to solve IVPs over an arbitrary time interval is discussed. Several modifications that improve the performance of MCPI methods for solving IVPs are presented. Simulation results of the applications for a small perturbation problem, a two-body problem, and a three-body problem are presented and compared with a conventional Runge-Kutta 4-5 integration method. Chapter IV focuses on using the MCPI methodology to solve BVPs. The strategy for solving different types of BVPs is illustrated using a linear second order system. The applications of the MCPI approach to the classical Lambert's two-point boundary value problem and an optimal trajectory design problem are also discussed in detail. Different techniques for enlarging the MCPI convergence domain are presented in Chapter V. The benefits of using these techniques are shown

through several examples. Chapter VI presents a parallel implementation of MCPI methods using a GPU card. In Chapter VII, conclusions and directions for future extension are discussed.

A. Existing Research in Solving IVPs and BVPs using Methods Related to MCPI Methods

1. Parallel Algorithms for Solving IVPs and BVPs

Compared with the significant achievement of using parallel computation techniques in other scientific computation fields, research on developing parallel algorithms to solve IVPs is advancing at a slower speed. This is because most of the current popular numerical integration methods do not have properties that lend themselves to parallel computation [1]. Franklin [2] compared three approaches to parallelize the existing forward integration methods: a parallel block implicit method, segmenting the equations to separate parts which can be solved using various processes, and revising the forward integration methods to a parallel-corrector form which was designed by Miranker and Liniger [3]. As many existing methodologies are limited to small scale problems (the number of the processors is small), Gear derived parallel algorithms targeted on large scale problems, but Gear's algorithms are only applicable for some special cases [1]. Although no illustrations were provided, Gear speculated that realizing a parallel structure for quadrature and function evaluation should be emphasized when designing more powerful methods. Gear's remarks provide a partial motivation and a qualitative context for this dissertation.

There has also been significant interest in using the parallel architectures for numerical optimization. Much of the recent progress can be categorized as either devising parallel linear algebra software and algorithms or developing new parallel

strategies in global optimization [4]. In the optimal control field, Travassos [5] suggested using Miranker and Liniger’s algorithms for parallel integration [3]. He also compared two parallel algorithms for function minimization, and discussed using a parallel shooting method to solve for the unknown initial costates. The level of parallelism of his approach is $n(2N - 1)$, where n is the dimension of the states and N is the number of subintervals used in the parallel shooting. Betts and Huffman [6] broke the overall trajectory into phases to parallelize the trajectory optimization algorithms. The sparse Jacobian matrix resulting from the multiple shooting method was solved using sparse finite differences by parallel computation. They implemented the algorithm on BBN GP100 (Butterfly) parallel processor for four orbit transfer problems and the speedups were in the range of 1.9 to 9.9. However, the results showed that their proposed parallel approach was not always faster than the serial direct approach, and the authors suggested this is because additional variables and constraints were introduced by the way they parallelized their problems.

2. Picard Iteration

Picard iteration is a successive solution approximation technique that is often used to prove the existence and uniqueness of the solutions to IVPs (detailed description in Appendix A). Although this approach is most often connected with the names of Charles Émile Picard, Ernst Lindelöf, Rudolph Lipschitz and Augustin Cauchy, it was first published by the French mathematician Joseph Liouville in 1838 for solving second order homogeneous linear equations. About fifty years later, Picard developed a much more generalized form, which placed the concept on a more rigorous mathematical foundation, and used this approach to solve boundary value problems described by second order ordinary differential equations [7]. The type of the boundary value problems he considered belongs to BVPs of the first kind (see definition in

Chapter IV). Picard developed a theoretical condition for the length of the intervals under which Picard iteration is guaranteed to converge. Including Picard himself, many authors have worked on improving these conditions [8, 9, 10, 11, 12]; generally these developments have sought to make Picard’s original convergence conditions sharper and less conservative. However, these conditions, even after several improvements, still only provide a conservative estimate on the region of the convergence; the “best possible” interval where Picard iteration converges, for a general problem, is not yet known. Also, there are a number of counterexamples in which Picard iteration is guaranteed not to converge, even if the starting function is arbitrarily close to the real solution [13, 14]. For multiple-point BVPs of the first kind, Urabe proved the sufficient conditions under which the approximate solution ensures the existence of the exact solution and presented an approach to calculate the error boundary for the approximation [15]. Coles and Sherman [12] are among the few who studied the convergence domain of Picard iteration for BVPs of the second and third kind (see definition in Chapter IV). Shridharana and Agarwal discussed how to construct a Picard-iteration-like sequence that will converge to the exact solutions [16]. However, except for some simple cases, their constructive theory is difficult to implement. Parker and Sochacki have studied the use of Picard iteration to generate solutions of IVPs in the form of Taylor series [17], however, convergence of this approach is not generally attractive.

3. Chebyshev Polynomials

Chebyshev polynomials are a complete set of orthogonal polynomials that are very important for function approximation (description in Appendix B). If the zeros of Chebyshev polynomials are used as the nodes for polynomial interpolations, the resulting approximating polynomial minimizes the Runge’s phenomenon and provides

the best approximation under the minimax norm [18]. Furthermore, for this set of nodes, the Chebyshev polynomials are orthogonal with a unit weight function.

Many researchers have contributed to the research on using Chebyshev polynomials to solve IVPs and BVPs [18, 19, 20, 21, 22, 23, 24, 25]. Urabe has made several important theoretical contributions. Using Newton's method and successive Galerkin's approximation, Urabe proved that the existence of an isolated periodic solution lying in the interior of the region where the differential equation is defined always implies the existence of Galerkin approximations of orders sufficiently high [23]. Furthermore, he showed that if there exists a Galerkin approximation of sufficiently high order, under some smoothness conditions, an exact solution exists and the approximation error can be determined. Urabe also discussed a numerical computation approach for periodic nonlinear systems using the Galerkin approximations [24]. Recently Chen presented the error bounds of the truncation errors when the two-variable Chebyshev series expansions are used to approximate the solutions of IVPs [25]. Urabe has also studied the use of Chebyshev polynomials to solve for the numerical solution of BVPs of the first kind [26]. He proved that if an isolated solution exists for multi-point BVPs, there are always Chebyshev approximate solutions that can be as accurate as desired. Furthermore, he showed that under some conditions, the obtained Chebyshev approximate solution insures the existence of the exact solution, and the error bound can be obtained. However, practical numerical methods for finding the Chebyshev coefficients using this approach are difficult to design (this assessment agrees with the discussion by Vlassenbroeck and Dooren [27]).

There have also been studies of using Chebyshev polynomials to solve optimal control problems, most of which belong to the set of *direct* methods. Vlassenbroeck and Dooren proposed using Chebyshev polynomials to approximate the state and control variables [27]. By approximating the performance function, dynamic equations,

and constraints, the original nonlinear optimal control problem is transformed into an algebraic equation system for the Chebyshev coefficients that are solved with nonlinear programming methods. Similar ideas have been used in other papers [28, 29], with the major differences of how the performance function and constraints are enforced.

4. Chebyshev-Picard Methods

Except for some special cases, it is usually difficult to use the classical Picard iteration method for solving IVPs and BVPs, because the integrals are not analytically tractable. Clenshaw and Norton [30] first proposed to solve IVPs and BVPs using both Picard iteration and Chebyshev polynomials (Chebyshev-Picard methods). They approximated the trajectory and the integrand by the same orthogonal basis functions (discrete Chebyshev polynomials) and integrated the basis functions term-by-term to establish a recursive trajectory approximation technique. Feagin applied the technique to several BVPs of the first kind [31]. Shave [32] studied using Chebyshev-Picard methods for orbit propagation and estimation problems based on the assumption of a single instruction, multiple data (SIMD) parallel architecture. He achieved high accuracy solutions for satellite motion propagation problems and also examined the potential time performance improvement that can be achieved by parallel implementation of Chebyshev-Picard methods. Recently, Sinha and Butcher developed a method that uses Picard iteration and shifted Chebyshev polynomials to symbolically solve for the approximate solutions of the state transition matrix for linear time-periodic dynamic systems [33]. In addition to the work by Shave [32], the parallel nature of the Chebyshev-Picard methods has also been addressed by Feagin and Fukushima. Feagin [34] presented a vector-matrix form of the Chebyshev-Picard method that is closely related to MCPI methods we propose in this dissertation. However, Feagin only applied the vector-matrix form of the Chebyshev-Picard methods

to solve IVPs. The capability to solve both IVPs and BVPs using this vector-matrix form was not discussed and he also did not provide any practical implementation or experimental results for solving either IVPs or BVPs. Fukushima implemented a Chebyshev-Picard algorithm on a vector computer [35]. However, for one example problem, the vector code was shown to be slower than the scalar code and the author suggested that this was because his approach could not be vectorized efficiently and the compilation put more additional overhead.

The issue of limited convergence domain for Chebyshev-Picard methods has been widely recognized since the beginning of this avenue of research, whereas the problem has not been finally solved up to today. In their pioneering paper, Clenshaw and Norton correctly concluded that the convergence for the proposed Chebyshev-Picard method, which is quite good in many problems, is not generally guaranteed. They suggested additional research to design more powerful methods for two-point BVPs [30]. Later, Norton proposed using what he called “Newton iteration formula” to solve the nonlinear ordinary differential equations in Chebyshev series [22]. The formulas for scalar ordinary differential equation are derived. He also discussed using iterative solutions to solve for the Chebyshev coefficients associated with the resulting linear equations for some special cases. Wright compared a Picard method, a linearization method, and some Taylor-series-based techniques for solving the problems associated with the Chebyshev collocation methods [21]. What he called the linearization method is essentially the same as the Newton iteration formula proposed by Norton. Wright showed several examples where the linearization method found the solutions but the classical Chebyshev-Picard method diverged. In his dissertation, Feagin extended the Newton iteration formula in Norton’s paper to vector cases and derived formulas for second order differential equations [31]. Realizing the difficulty of solving the resulting linear equations for high dimensional problems using the linearization

approach, Feagin discussed two methods for improving the direct inversion approach used by Norton: one is to solve the linear equations using the iterative method proposed by Scraton [20]; the other is to use the successive over-relaxation method [36], which he showed to be the most efficient.

5. Features of the Modified Chebyshev-Picard Iteration Methods

Motivated by the above research, this dissertation introduces a family of modified Chebyshev-Picard iteration (MCPI) methods. The proposed MCPI methods, developed in subsequent chapters, have following four unique features:

- This is the first time that a unified vector-matrix form of Chebyshev-Picard methods is developed and proven to be applicable to solve both IVPs and BVPs.
- This is the first time that the power of Chebyshev-Picard methods for solving both IVPs and BVPs has been shown through highly accurate and efficient solution of several important celestial mechanics problems (even prior to parallel implementation).
- Several new techniques are proposed in this dissertation that aim to either improve the accuracy performance or enlarge the convergence domain for Chebyshev-Picard methods. This is the first time that the perturbation motion MCPI methods are proposed and also shown to improve the performance. The concept of the correction control in Picard iterations is novel. Unlike other existing techniques to enlarge the convergence domain, this method does not encounter the curse of dimensionality difficulties when solving high dimensional problems or when high order polynomials are required to approximate the solutions.
- This is the first time that Chebyshev-Picard methods have been implemented

on a graphics card to obtain a parallel implementation. The speedup achieved from the parallel implementation is the largest that has ever been reported.

We point out that this dissertation has a strong connection with Feagin's work [31, 34]. In his dissertation, Feagin proposed to combine Chebyshev polynomials with either Picard iterations or quasilinearization to solve two-point BVPs for spacecraft trajectory design. He also recommended the utilization of Aitken's process to enlarge the convergence domain as we do in this dissertation. In the appendix of his dissertation, Feagin derived a vector-matrix form of Chebyshev-Picard methods for solving IVPs, which is very similar to the vector-matrix form of MCPI methods we introduce in Chapter II. However, the differences of this dissertation from the work of Feagin are significant. In addition to the four distinguishing features we summarized earlier, we emphasize the following five differences between Feagin's work and the results documented in this dissertation:

- The vector-matrix form of MCPI methods we propose is different from the form Feagin derived. The differences are pointed out in detail in Chapter II.
- The applicability of the vector-matrix form of MCPI methods for solving BVPs was neither discussed nor used by Feagin.
- Feagin focused on solving the BVPs of the first type (see definition in Chapter IV), whereas we discuss using MCPI methods to solve more general BVPs.
- Feagin did not show any example about using the Chebyshev-Picard method for solving IVPs.
- Feagin did consider in his dissertation an optimal control problem, however, in lieu of applying a Picard iteration, he instead used a quasilinearization approach.

In this dissertation we discuss using MCPI methods to solve general family of optimal control problems and we prove the practical merit of this approach.

We mention that the above differences are pointed out for clarity in understanding the context of this dissertation's contribution, and certainly in no way minimize the pioneering and important contributions made by Feagin, whose work proceeded the present work by almost four decades!

B. Application Areas

We will study the following IVP and BVP applications that are very important in the celestial mechanism fields.

1. Orbit Propagation

Since the launch of the first satellite (Sputnik 1) in 1957, it has been necessary to maintain and frequently update database of satellite orbits. This endeavor is referred to maintaining the space catalog. The ability to propagate satellite motion quickly and accurately is one of the major factors that affect the completeness, accuracy, and cost of the database [37], especially as the number of objects in the catalog exceeds 100,000. The analytic predictor methods, while very efficient, can not meet requirements for the task such as collision avoidance and satellite acquisition. Therefore numerical integration of the satellite motion with even more accurate and complicated perturbation force models has become necessary [38]. The most widely accepted integrator is an 8th-order multi-step predictor-corrector method known as the Gauss-Jackson method [38, 39, 40]. However, this choice of integrator does not allow the SIMD computer to run very efficiently [39]. With the current goal to model around 150,000 satellites within the next 8 to 10 years with a degree and order 36×36 or

higher gravity model and a realistic atmospheric density model, and a requirement to estimate each of these orbits in near-real-time as observations become available, the extensive force calculation has become a processing bottleneck. Even with the anticipated speedup from Moore's Law, a true parallel algorithm has become necessary.¹

2. Lambert's Problem

Lambert's problem is the classical two-point BVP of celestial mechanics that was first stated and solved by Johann Heinrich Lambert in 1761 [41]. Given two positions at an initial time and a prescribed final time, Lambert's problem is to solve for an initial velocity, using which the generated orbit, connects the two known positions with the prescribed time of flight. For general perturbed motion, the generalized Lambert's problem remains a challenging nonlinear problem that must be frequently solved today for orbit transfer. Battin developed an elegant, now classical approach to solve the classical Lambert's problem [41]. However, his method only works for the cases that no perturbations to the inverse-square gravity field exist. Numerical iteration techniques such as shooting methods [42] are the most frequently used current approach to solve the more general perturbed orbit transfer problems.

3. Optimal Control Problems

Given a performance function, a set of ordinary differential equations describing the dynamic system, and various system constraints, an optimal control design algorithm solves for the control input that drives the system to execute an optimal trajectory to minimize or maximize the performance function and satisfy the constraints. The two most popular sets of computational techniques for solving optimal control problems

¹Communication with Dr. Paul W. Schumacher from HSAI-SSA, Air Force Research Laboratory

are *direct* shooting and *indirect* shooting[43] methods. *Direct methods* introduce a parametric representation of the control variables (and frequently the state variables as well), and then use nonlinear programming optimizers to solve the resulting nonlinear parameter optimization problems (a so-called nonlinear programming problem). With the increasing power of these optimizers, direct approaches can obtain sub-optimal solutions often more easily than indirect approaches. However, the accuracy and optimality of the returned control policy for the original continuous system are not guaranteed. *Indirect approaches* are based on calculus of variations without parameterization of state and control variables. Necessary conditions are derived from Pontryagin's Principle[44]. A simple shooting or multiple shooting method is usually used to solve the resulting two-point value problem, with the goal being to find the unknown initial states and co-states. The solutions obtained from the indirect approaches assure local optimality and accuracy. However, solving the resulting BVPs is challenging because the solutions can be very sensitive to the initial guess, and local convergence to solutions that satisfy the necessary conditions does not guarantee global optimality.

Through Pontryagin's Principle, we consider the optimal control of continuous systems as a special type of BVPs in this dissertation. We introduce MCPI methods to solve the state and co-state differential equations, and boundary conditions that are the necessary conditions for optimality.

CHAPTER II

PROBLEM STATEMENT AND METHODOLOGY OF MCPI METHODS

A. Introduction

Formal statements of IVPs and BVPs that are studied in this dissertation are presented in this chapter. The MCPI methodology is discussed and its vector-matrix forms to solve IVPs and BVPs are developed. The formulations presented in this chapter provide the bases for the later chapters to solve IVPs, BVPs, and optimal control problems.

B. Problem Statement of IVPs and BVPs

The systems we consider in this dissertation are described by ordinary differential equations (ODEs), for which the independent variable is denoted by the time t and the vector of dependent variables is represented by \mathbf{x} , whose elements include all the states of the dynamic system. Here we assume higher order ODEs have already been transformed to their first order forms.

An IVP seeks the solution of $\mathbf{x}(t)$ that satisfies the given ODE

$$\frac{d\mathbf{x}(t)}{dt} = \mathbf{f}(t, \mathbf{x}(t)), t \in [a, b] \quad (2.1)$$

and the given initial condition $\mathbf{x}(t = a) = \mathbf{x}_0$, where $t = a$ is the initial time and $t = b$ is the final time. The conditions for the existence and uniqueness of the IVP solution are the same as the conditions for the classical Picard iterations to converge (presented in Appendix A).

Although general BVPs can have interior constraint conditions on \mathbf{x} at any time in the interval of $[a, b]$, we focus on two-point BVPs in this dissertation. In this case,

a BVP seeks the solution for $\mathbf{x}(t)$ that satisfies the given ODE in Eq. 2.1 and the boundary conditions for \mathbf{x} , some of which are defined at the initial time and others are defined at the final time.

C. Modified Chebyshev-Picard Iteration Methods

We use a scalar case example to illustrate the procedure to solve IVPs using MCPI methods. Consider a differential equation with an initial condition $x(t = a) = x_0$

$$\frac{dx}{dt} = f(t, x), t \in [a, b] \quad (2.2)$$

The first step of MCPI methods is to transform the generic independent variable t to a new variable τ , which is defined on the valid range $-1 \leq \tau \leq 1$ of Chebyshev polynomials through

$$t = \frac{b-a}{2}\tau + \frac{b+a}{2} \quad (2.3)$$

After the substitution, Eq. (2.2) is transformed to a new form as

$$\frac{dx}{d\tau} = \frac{b-a}{2} f\left(\frac{b-a}{2}\tau + \frac{b+a}{2}, x\right) \equiv g(\tau, x) \quad (2.4)$$

The Chebyshev polynomial of degree k is denoted by T_k and the $(N+1)$ discrete nodes that are used to approximate the states are the Chebyshev-Gauss-Lobatto (CGL) nodes, which are calculated from

$$\tau_j = \cos(j\pi/N), j = 0, 1, \dots, N \quad (2.5)$$

To start the iteration, an initial guess of the solution $x^0(\tau)$ is required. Assume the force function appearing on the right hand side of Eq. 2.4 is approximated by a

N^{th} order Chebyshev polynomial

$$\frac{dx}{d\tau} = g(\tau, x^0) \approx \sum_{k=0}^{k=N} F_k T_k(\tau) \quad (2.6)$$

Using the discrete orthogonality property of Chebyshev polynomials, the coefficients $\{F_0, F_1, \dots, F_N\}$ are calculated immediately from

$$F_k = \frac{2}{N} \sum_{j=0}^N g(\tau_j, x^0(\tau_j)) T_k(\tau_j) \quad (2.7)$$

Notice each coefficient F_k is obtained through the summation of $(N + 1)$ independent terms, each of which is the multiplication of the force function g and the Chebyshev polynomials T_k evaluated at the CGL point τ_j . Furthermore, all the coefficients $\{F_0, F_1, \dots, F_N\}$ are independent of each other, and can therefore be computed in parallel processors. Also, for problems where calculating the force function g is time consuming, significant time performance improvement can be achieved by parallel computation of g on different processors.

The solution at the next step is assumed as $x^1(\tau) = \sum_{k=0}^{k=N} \beta_k T_k(\tau)$, and the Picard method provides the iteration form

$$x^1(\tau) = x_0 + \int_{-1}^{\tau} g(s, x^0(s)) ds \quad (2.8)$$

Using the integration properties of Chebyshev polynomials (see Appendix B), using Eq. 2.6 in Eq. 2.8, one obtains

$$\begin{aligned} & \frac{1}{2} \beta_0 T_0(\tau) + \beta_1 T_1(\tau) + \beta_2 T_2(\tau) + \dots + \beta_N T_N(\tau) \\ &= \mathbf{x}(-1) + \int_{-1}^{\tau} \left(\frac{1}{2} F_0 T_0(s) + F_1 T_1(s) \dots + F_N T_N(s) \right) ds \\ &= c + \frac{1}{2} F_0 T_1(\tau) + \frac{1}{4} F_1 T_2(\tau) + \frac{1}{2} F_2 \left[\frac{1}{3} T_3(\tau) - T_1(\tau) \right] \dots \\ & \quad + \frac{1}{2} F_N \left[\frac{1}{N+1} T_{N+1}(\tau) - \frac{1}{N-1} T_{N-1}(\tau) \right] \end{aligned} \quad (2.9)$$

where c is some constant to be defined according to the given constraints.

For the IVP with the given initial condition $x(t = a) = x_0$, equating the first to N^{th} order coefficients of the Chebyshev polynomials on the left side and on the right side of Eq. 2.9 leads to the following formulations to obtain the coefficients

$$\beta_r = \frac{1}{2r}(F_{k-1} - F_{k+1}), r = 1, 2, \dots, N - 1 \quad (2.10)$$

$$\beta_N = \frac{F_{N-1}}{2N} \quad (2.11)$$

The initial condition leads to the solution for the zero order coefficient, given by

$$\beta_0 = 2x_0 + 2\sum_{k=1}^{N} (-1)^{k+1} \beta_k \quad (2.12)$$

For two-point BVPs, both the zero and first order coefficients should be calculated to satisfy the boundary conditions $x(t = a) = x_0$ and $x(t = b) = x_f$, leading to the update equations for the state approximation coefficients as

$$\beta_r = \frac{1}{2r}(F_{k-1} - F_{k+1}), r = 2, 3, \dots, N - 1 \quad (2.13)$$

$$\beta_N = \frac{F_{N-1}}{2N} \quad (2.14)$$

$$\beta_0 = x_0 + x_f - 2(\beta_2 + \beta_4 + \beta_6 + \dots) \quad (2.15)$$

$$\beta_1 = \frac{(x_f - x_0)}{2} - (\beta_3 + \beta_5 + \beta_7 + \dots) \quad (2.16)$$

The updated coefficients are used to calculate the new trajectory approximation for the next step. In this way, the solutions are iteratively improved until some accuracy requirements are satisfied. To account for the nonlinearity issues, the stopping criterion we choose is to require both the difference between the solutions x^i and x^{i-1} and the difference between the solutions x^i and x^{i+1} are less than some tolerance. Figure 1 shows the steps of MCPI methods for solving IVPs and Fig. 2 shows the steps of MCPI methods for solving BVPs. Notice when solving BVPs, the MCPI

algorithm differs significantly from usual shooting methods: there is no local Taylor series approximations, gradient computations, or matrix inversions. Furthermore, as a consequence of using an accurate Chebyshev approximation of the integrand on each iteration, the integration of the Chebyshev solutions is accomplished analytically when the coefficients are updated.

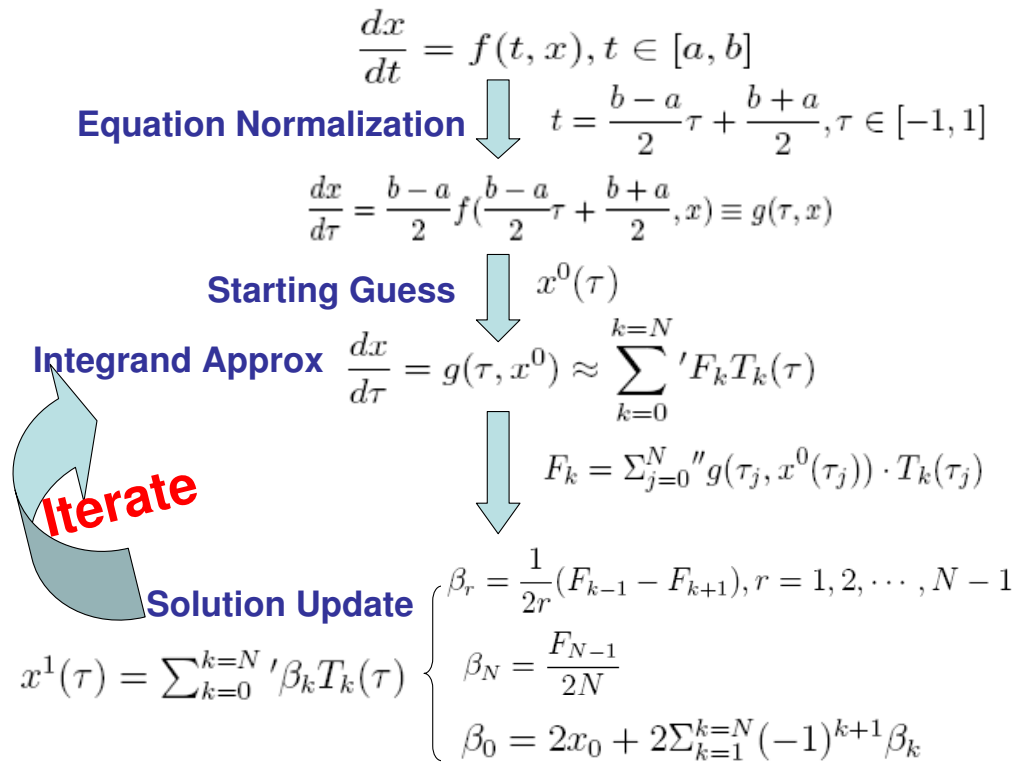


Fig. 1. Steps of MCPI methods for solving IVPs

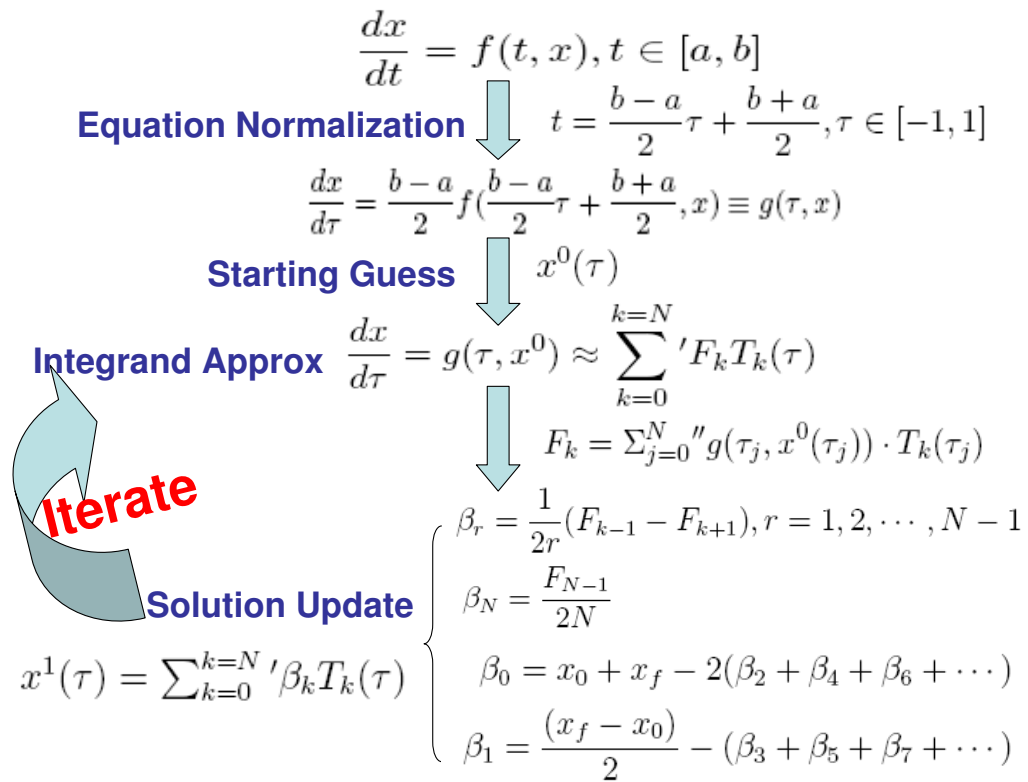


Fig. 2. Steps of MCPI methods for solving BVPs

D. Vector-Matrix Forms of MCPI Methods

Instead of term by term to solve for the state value at the $(N + 1)$ CGL nodes, the $(N + 1)$ Chebyshev coefficients, and the updated $(N + 1)$ Chebyshev coefficients, we introduce a compact vector-matrix approach to implement MCPI methods. In this way, we not only solve for the state values and Chebyshev coefficients in a vector form, but also has a compact formulation to update the coefficients through Picard iterations. We find that a similar idea has been published by Feagin and Nacozy[34]. There are three major differences between the formulations in their paper and our derivations. We will point out these differences along the way we derive for the vector-matrix form of MCPI methods.

The zero to N^{th} order coefficients of the Chebyshev polynomial to approximate the solution $x(\tau)$ are defined in a vector form as

$$\vec{\alpha} = [\alpha_0, \alpha_1, \dots, \alpha_N]^T \quad (2.17)$$

The solution of x evaluated at the CGL nodes is represented by a vector as

$$\vec{x} = [x(\tau_0), x(\tau_1), \dots, x(\tau_N)]^T \quad (2.18)$$

Similar as Eq. 2.6, the solution of x can be calculated from its Chebyshev coefficients

through

$$\begin{aligned}
\vec{x} &= \begin{bmatrix} \frac{1}{2}\alpha_0 T_0(\tau_0) + \alpha_1 T_1(\tau_0) + \cdots + \alpha_N T_N(\tau_0) \\ \frac{1}{2}\alpha_0 T_0(\tau_1) + \alpha_1 T_1(\tau_1) + \cdots + \alpha_N T_N(\tau_1) \\ \vdots \\ \frac{1}{2}\alpha_0 T_0(\tau_N) + \alpha_1 T_1(\tau_N) + \cdots + \alpha_N T_N(\tau_N) \end{bmatrix} \\
&= \begin{bmatrix} T_0(\tau_0) & T_1(\tau_0) & \cdots & T_N(\tau_0) \\ T_0(\tau_1) & T_1(\tau_1) & \cdots & T_N(\tau_1) \\ \vdots & \vdots & \vdots & \vdots \\ T_0(\tau_N) & T_1(\tau_N) & \cdots & T_N(\tau_N) \end{bmatrix} \begin{bmatrix} \frac{1}{2} & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & \ddots & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} \alpha_0 \\ \alpha_1 \\ \vdots \\ \alpha_N \end{bmatrix} \equiv TW\vec{\alpha} \\
&\equiv C_x\vec{\alpha} \quad (2.19)
\end{aligned}$$

with the definition

$$C_x \equiv TW \quad (2.20)$$

$$T = \begin{bmatrix} T_0(\tau_0) & T_1(\tau_0) & \cdots & T_N(\tau_0) \\ T_0(\tau_1) & T_1(\tau_1) & \cdots & T_N(\tau_1) \\ \vdots & \vdots & \vdots & \vdots \\ T_0(\tau_N) & T_1(\tau_N) & \cdots & T_N(\tau_N) \end{bmatrix} \quad (2.21)$$

and the diagonal matrix W is defined as

$$W = \text{diag}([\frac{1}{2}, 1, 1, \cdots, 1, 1]) \quad (2.22)$$

Notice the last diagonal term in Eq. 2.22 is one since we are using the discrete least-squares fit (see Appendix B). Although Feagin and Nacozy[34] did not say explicitly the reason for their formula to have the last term as $\frac{1}{2}$, that is required for exact interpolation where the function $f(x, t)$ fits exactly at the interpolation nodes. The most important reason we choose the least-square fit approach is because this approach implicitly provides an error bound while for the interpolation approach, this bound

will be difficult to determine.

The force function is evaluated at the CGL nodes and defined in a vector form as

$$\vec{g} = [g(\tau_0), g(\tau_1), \dots, g(\tau_N)]^T \quad (2.23)$$

The zero to N^{th} order coefficients of the Chebyshev polynomials for function g are defined as

$$\vec{F} = [F_0, F_1, \dots, F_N]^T \quad (2.24)$$

Consistently with Eq. 2.7, these coefficients are calculated from

$$\begin{aligned} \vec{F} &= \begin{bmatrix} \frac{1}{N}g(\tau_0)T_0(\tau_0) + \frac{2}{N}g(\tau_1)T_0(\tau_1) + \dots + \frac{1}{N}g(\tau_N)T_0(\tau_N) \\ \frac{1}{N}g(\tau_0)T_1(\tau_0) + \frac{2}{N}g(\tau_1)T_1(\tau_1) + \dots + \frac{1}{N}g(\tau_N)T_1(\tau_N) \\ \vdots \\ \frac{1}{N}g(\tau_0)T_N(\tau_0) + \frac{2}{N}g(\tau_1)T_N(\tau_1) + \dots + \frac{1}{N}g(\tau_N)T_N(\tau_N) \end{bmatrix} \\ &= \begin{bmatrix} T_0(\tau_0) & T_0(\tau_1) & \dots & T_0(\tau_N) \\ T_1(\tau_0) & T_1(\tau_1) & \dots & T_1(\tau_N) \\ \vdots & \vdots & \vdots & \vdots \\ T_N(\tau_0) & T_N(\tau_1) & \dots & T_N(\tau_N) \end{bmatrix} \begin{bmatrix} \frac{1}{N} & 0 & 0 & 0 \\ 0 & \frac{2}{N} & 0 & 0 \\ 0 & 0 & \dots & 0 \\ 0 & 0 & 0 & \frac{1}{N} \end{bmatrix} \begin{bmatrix} F_0 \\ F_1 \\ \vdots \\ F_N \end{bmatrix} \\ &\equiv T^T V \vec{g} \\ &= TV \vec{g} \quad (2.25) \end{aligned}$$

where the property that $T^T = T$ has been utilized. And the diagonal matrix V is defined as

$$V = \text{diag}\left(\left[\frac{1}{N}, \frac{2}{N}, \frac{2}{N}, \dots, \frac{2}{N}, \frac{1}{N}\right]\right) \quad (2.26)$$

The zero to N^{th} order coefficients of the Chebyshev polynomials to approximate

the updated solution of $x(\tau)$ are defined as

$$\vec{\beta} = [\beta_0, \beta_1, \dots, \beta_N]^T \quad (2.27)$$

The formulations from Eqs. (2.10) to (2.12) for updating the coefficients correspond

to a vector-matrix form as

$$\begin{aligned}
\vec{\beta} &= \begin{bmatrix} 2x_0 + 2(\beta_1 - \beta_2 + \beta_3 + \cdots + (-1)^{N+1}\beta_N) \\ \frac{1}{2}(F_0 - F_2) \\ \frac{1}{2 \times 2}(F_1 - F_3) \\ \vdots \\ \frac{1}{2 \times r}(F_{r-1} - F_{r+1}) \\ \vdots \\ \frac{F_{N-1}}{2N} \end{bmatrix} \\
&= \begin{bmatrix} 2x_0 \\ 0 \\ 0 \\ \vdots \\ 0 \\ 0 \end{bmatrix} \\
&+ \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & \frac{1}{2} & 0 & 0 & 0 \\ 0 & 0 & \frac{1}{4} & 0 & 0 \\ 0 & 0 & 0 & \ddots & 0 \\ 0 & 0 & 0 & \frac{1}{2 \times r} & 0 \\ 0 & 0 & 0 & \ddots & 0 \\ 0 & 0 & 0 & 0 & \frac{1}{2N} \end{bmatrix} \begin{bmatrix} 1 & -\frac{1}{2} & -\frac{2}{3} & \frac{1}{4} & -\frac{2}{15} & \cdots & (-1)^{N+1} \frac{1}{N-1} \\ 1 & 0 & -1 & 0 & 0 & \cdots & 0 \\ 0 & 1 & 0 & -1 & 0 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & \cdots & 1 & 0 & -1 \\ 0 & 0 & 0 & 0 & \cdots & 1 & 0 \end{bmatrix} \begin{bmatrix} F_0 \\ F_1 \\ F_2 \\ F_3 \\ \vdots \\ F_{N-1} \\ F_N \end{bmatrix} \\
&= \vec{\chi}_0 + RS\vec{F}
\end{aligned}$$

which can be further written as

$$\vec{\beta} = \vec{\chi}_0 + RSTV\vec{g} = \vec{\chi}_0 + C_\alpha\vec{g} \quad (2.28)$$

as we define

$$RSTV \equiv C_\alpha \quad (2.29)$$

$$\vec{\chi}_0 = [2x_0, 0, 0, \dots, 0]^T \quad (2.30)$$

$$R = \text{diag}\left([1, \frac{1}{2}, \frac{1}{4}, \dots, \frac{1}{2(N-1)}, \frac{1}{2N}]\right) \quad (2.31)$$

$$S = \begin{bmatrix} 1 & -\frac{1}{2} & -\frac{2}{3} & \frac{1}{4} & -\frac{2}{15} & \dots & (-1)^{N+1} \frac{1}{N-1} \\ 1 & 0 & -1 & 0 & 0 & \dots & 0 \\ 0 & 1 & 0 & -1 & 0 & \dots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & \dots & 1 & 0 & -1 \\ 0 & 0 & 0 & 0 & \dots & 1 & 0 \end{bmatrix} \quad (2.32)$$

The first row of matrix S is obtained as we represent β_1 to β_N in terms of F_0 to F_N .

The r^{th} ($r = 2, 3, \dots, N-1$) column of this row has a form as

$$S[1, r] = (-1)^{r+1} \left(\frac{1}{r-1} - \frac{1}{r+1} \right) \quad (2.33)$$

We think the vector form of R defined by Feagin and Nacozy[34] may contain a typographical error. Additionally, Feagin and Nacozy introduced a new dependent variable to transform the general initial condition to zero initial condition, which is the third difference between their approach and our approach. The compact form of Eq. 2.28 helps to extend MCPI methods to solve BVPs, which Feagin and Nacozy did not discuss. For BVPs, the zero and first order coefficients are calculated from Eqs. 2.15 and 2.16. This vector-matrix form of MCPI methods is computationally more efficient than using a “for loop” computation to recursively evaluate the original scalar form. Additionally, since both C_x and C_α are constant once the order of polynomials is fixed, these matrices can be computed once before the iteration starts, which results in significant speedup for the problems that require high order poly-

mials to approximate solutions. A flowchart to implement the vector-matrix form of MCPI methods for solving IVPs is shown in Fig 3.

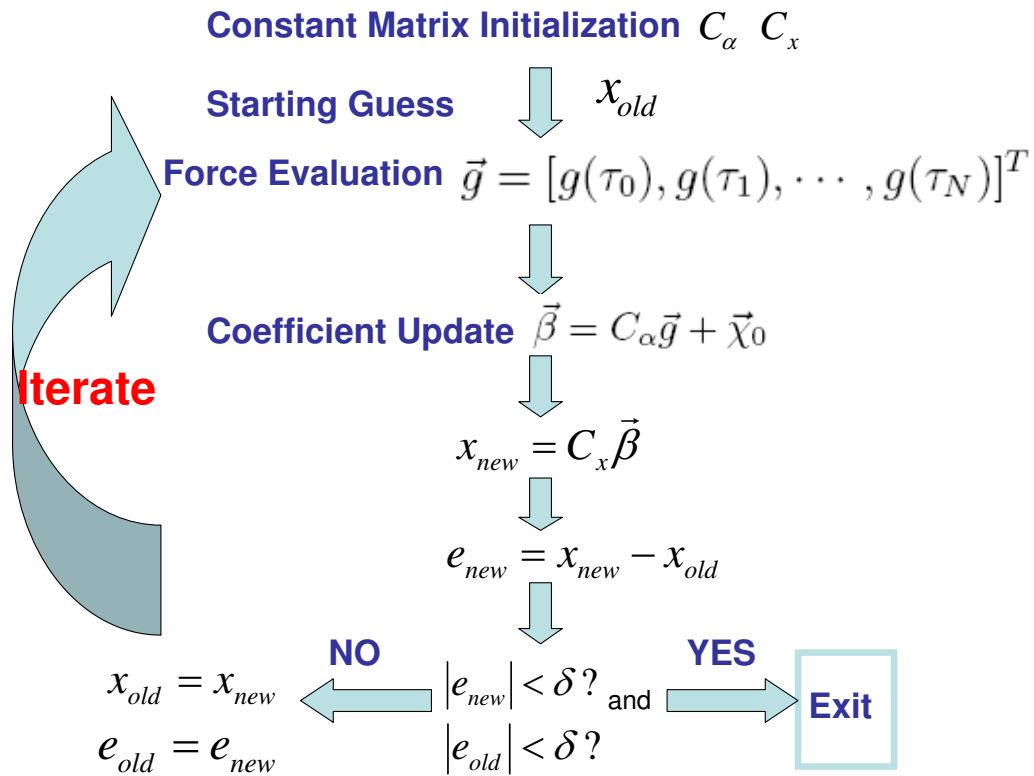


Fig. 3. Flowchart of the vector-matrix form of the MCPI Approach

E. Summary

This chapter develops a unified framework of MCPI methods that can be used for solving IVPs and BVPs. Compared with most existing methods that are used for solving IVPs and BVPs, the proposed methods are inherently parallel algorithms. Significant speedup can be achieved through large scale parallel computation, especially for the cases where the force function evaluation is computationally expensive. The solutions obtained from MCPI methods are in the form of Chebyshev polynomials, thus the interpolated values of the states at any time can be obtained immediately. We use the proposed MCPI methods to solve IVPs in the next chapter.

CHAPTER III

MODIFIED CHEBYSHEV-PICARD ITERATION METHODS FOR SOLUTION
OF INITIAL VALUE PROBLEMS

A. Introduction

Using the MCPI methodology developed in Chapter II, this chapter focuses on using MCPI methods for solving IVPs. The chapter starts by analyzing the convergence characteristics of MCPI methods through a linear scalar case example. Analogous to analytical continuation, we then introduces a piecewise approach which extends MCPI methods to be able to solve IVPs on arbitrarily long time intervals, and for many problems effectively guarantee that accurate converged solutions can be attained. Several techniques to improve the performance of MCPI methods are presented, which include approximating only the position coordinates for second order systems and integrating only the perturbation motion based on a reference motion. The power of the MCPI approach is illustrated in the numerical example section through studying a nonlinear perturbation from the sinusoidal motion problem, then through solution of more practical problems of celestial mechanics: two-body problem with and without zonal harmonic gravitational perturbations, and a three-body problem. After discussing these examples, we summarize insight on how to tune several parameters for MCPI methods.

B. Convergence Analysis of MCPI Methods

Because of the accumulation of round off and approximation errors during the iterations (when a finite order of Chebyshev polynomial is used to approximate solutions), the convergence domain of MCPI methods is different from the ideal conditions under

which Picard iteration theoretically converges (see Appendix A). This is qualitatively similar to all the conventional single-step and multi-step numerical methods for solving nonlinear differential equations, which have their own convergence limitations (which typically require automatic step size control, and occasionally, artistic tuning to achieve convergence). Analogously, establishing the rigorous convergence domain of MCPI methods applicable for general nonlinear systems is not possible by any known approach. Instead, we first use a linear scalar problem as an example to show that the global convergence of MCPI methods is not generally guaranteed, and we then address the practical issue of checking the convergence, and importantly, approaches to enlarge the convergence domain.

Consider a linear dynamic system

$$\frac{dx}{dt} = cx(t), t \in [a, b] \quad (3.1)$$

with an initial condition $x(t = a) = x_0$. First the generic independent variable t is transformed linearly to the range $[-1, 1]$ through

$$t = \frac{b-a}{2}\tau + \frac{b+a}{2}, \tau \in [-1, 1] \quad (3.2)$$

After this time normalization, using $x(t) = X(\tau)$, Eq. 3.1 is transformed to

$$\frac{dX}{d\tau} = \frac{b-a}{2}cX(\tau) \quad (3.3)$$

The k^{th} step solution of X evaluated at the $N + 1$ CGL nodes is represented by a vector

$$\vec{X}_k = [X(\tau_0), X(\tau_1), \dots, X(\tau_N)]^T \quad (3.4)$$

An initial condition vector Θ_0 is defined as

$$\Theta_0 = [2x_0, 0, 0, \dots, 0]^T, \Theta_0 \in R^{N+1} \quad (3.5)$$

Consistent with Eq. 2.28, $\vec{\beta}$, which is the coefficient vector of the Chebyshev polynomials to approximate the solution of \vec{X} , is obtained through

$$\vec{\beta} = C_\alpha \frac{b-a}{2} c \vec{X}_k + \Theta_0 \quad (3.6)$$

Using Eq. 2.19, Picard iteration leads to the solution at the next step as

$$\vec{X}_{k+1} = C_x (C_\alpha \frac{b-a}{2} c \vec{X}_k + \Theta_0) \quad (3.7)$$

The solutions are iteratively updated until the stopping criterion is satisfied. Defining two constant matrices K_1 and K_2 as

$$K_1 \equiv \frac{b-a}{2} c C_x C_\alpha \quad (3.8)$$

$$K_2 \equiv C_x \Theta_0 \quad (3.9)$$

we obtain a compact form of the MCPI method for this linear case as

$$\vec{X}_{k+1} = K_1 \vec{X}_k + K_2, k = 1, 2, \dots, \quad (3.10)$$

It is known in the linear system theory that the sequence in Eq. 3.10 is convergent only if all the eigenvalues of K_1 are within a unit circle. Equation 3.8 shows that the convergence of the MCPI method is dependent on both the dynamical system characteristics “ c ” and the length of the time interval ($b - a$). A long interval can lead to the divergence of the method while reducing the time interval can improve the convergence of the method. The maximum eigenvalue of $C_x C_\alpha$ is denoted as $\lambda_{max}(C_x C_\alpha)$, which dictates convergence of Eq. 3.10 (see Eq. 3.8). In Figs. 4 to 7, we plot $\lambda_{max}(C_x)$, $\lambda_{max}(C_\alpha)$, and $\lambda_{max}(C_x C_\alpha)$. Note $\lambda_{max}(C_x)$ vs N is approximately proportional to \sqrt{N} , whereas $\lambda_{max}(C_\alpha)$ decays rapidly vs N . $\lambda_{max}(C_x C_\alpha)$ has more complicated behavior as is evident in Figs. 6 and 7. For small N ($N < 40$), $\lambda_{max}(C_x C_\alpha)$

decays from about 0.7 to about 0.05, almost linearly on a log-log scale. Thereafter, for $N > 40$, $\lambda_{max}(C_x C_\alpha) \approx 0.05$, remaining approximately constant. As is evident in Fig. 7, this trend holds for large N . Specially, referring to Eq. 3.8, we require $\lambda_{max}(K_1) = \frac{b-a}{2} c \lambda_{max}(C_x C_\alpha)$ to be in the unit circle. This gives rise to the maximum interval length

$$(b - a)_{max} = \frac{2}{c} \frac{1}{\lambda_{max}(C_x C_\alpha)} \quad (3.11)$$

In this case, since $\lambda_{max}(C_x C_\alpha) \approx 0.05$ for $N > 40$, we have

$$(b - a)_{max} = \frac{40}{c} \quad (3.12)$$

if c is in the unit of $\frac{1}{\lambda_{max}(C_x C_\alpha)}$. While $\lambda_{max}(K_1) < 1$ guarantees convergence of the Picard iterations, for a fixed N , it does not guarantee that N is sufficiently high. It is fortunate, as is evident in Fig. 7, that convergence does not degrade for large N .

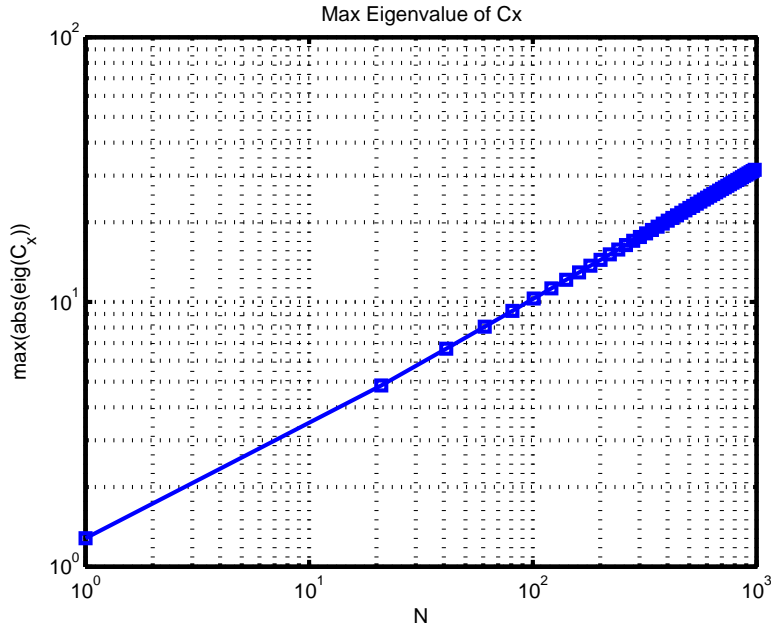
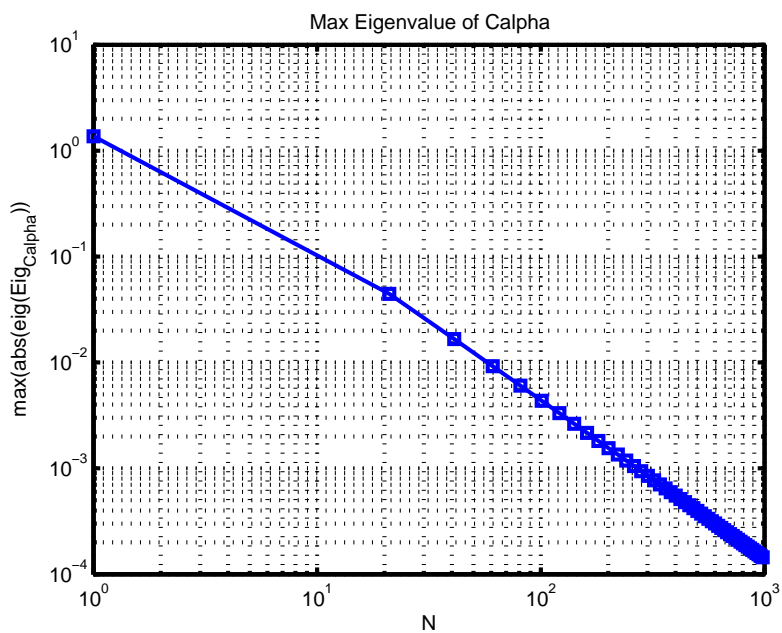
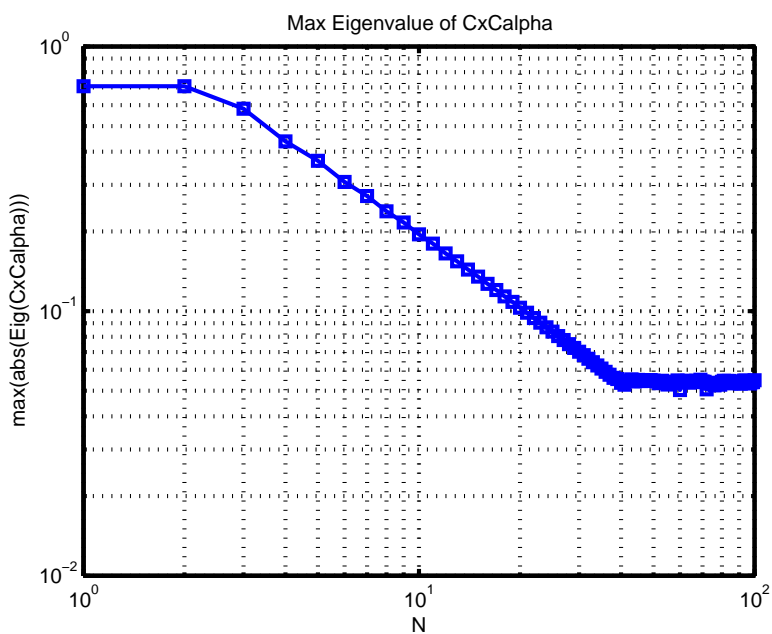


Fig. 4. Max eigenvalue of C_x

Fig. 5. Max eigenvalue of C_{α} Fig. 6. Max eigenvalue of $C_x C_{\alpha}$ (N:1-100)

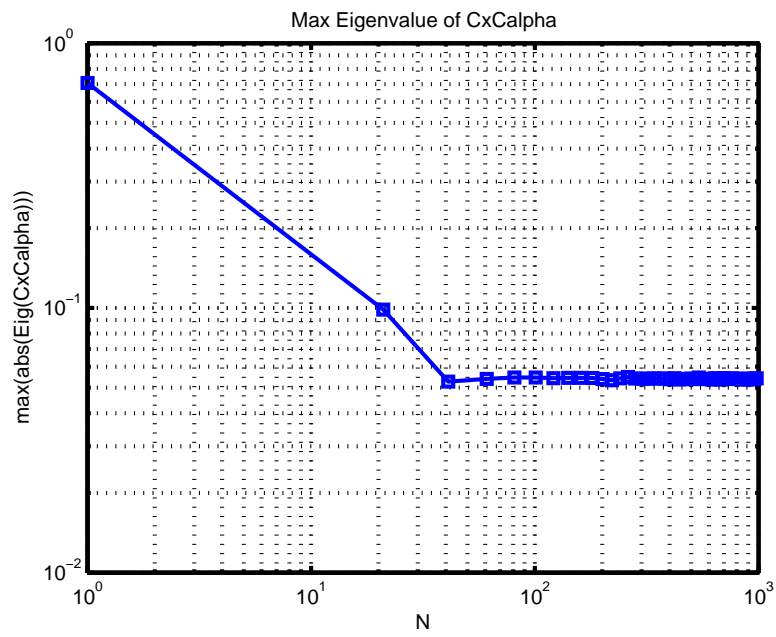


Fig. 7. Max eigenvalue of $C_x C_\alpha (N : 1 - 1000)$

C. Piecewise MCPI Methods for Solving IVPs

Although the Chebyshev-Picard iteration algorithm only converges on a finite interval, we can anticipate using a piecewise approach to solve a significant family of IVPs over a large domain. The initial conditions on the next segment are the final state values on the previous segment. This may sound similar to the concept of the step size control used in forward integration methods such as Runge-Kutta methods. However, the step size used by MCPI methods is typically a much larger finite interval than the possible step sizes used by the typical numerical methods, as will be shown in the numerical examples. Furthermore, compared with the forward integration methods in which the integration errors are typically increasing with time in a weakly unstable fashion, better accuracy can be achieved from using MCPI methods because the largest errors from MCPI methods usually appear in the middle of the interval and the smallest errors are at the ends where adjacent (successive) segments are joined. The fundamental reason for this special characteristic of MCPI methods is because of the chosen Chebyshev basis functions and CGL nodes which are denser at the boundaries and sparser in the middle.

D. Using MCPI Methods to Solve Second Order IVPs with Only Position Integration

Although we can always transform a system of second order ODEs to a system of first order ODEs, MCPI methods can be further simplified such that we can apply the Chebyshev approximation only to capture the position and acceleration states. Physically, the derivative of the position is always the velocity, thus we theoretically do not obtain additional information when we integrate the position to get the velocity during the process of iterations. Note the distinction, if we abandon the double integration of acceleration to position, in favor of doubling the dimension and applying

the first order version of Picard iterations, we end up writing independent Chebyshev approximations for position and velocity, and are led to twice as many coefficients. Computationally, by integrating twice the acceleration to update only the position, we reduce the dimension of the problem by half so some speedup and dimension reduction can be obtained by not solving for the velocity. Through the derivative properties of Chebyshev polynomials, the velocity approximation can be obtained immediately from the solutions for the position.

Consider a scalar case second order dynamic system

$$\ddot{x} = f(x(t), t), t \in [a, b] \quad (3.13)$$

the initial position is $x(t = a) = x_0$ and the initial velocity is $\dot{x}(t = a) = \dot{x}_0$. The generic independent variable t is translated to the range $[-1, 1]$ through

$$t = \frac{b-a}{2}\tau + \frac{b+a}{2}, \tau \in [-1, 1] \quad (3.14)$$

After this time normalization, Eq. 3.13 is transformed to

$$\frac{dx}{d\tau} = \frac{b-a}{2}\dot{x}(\tau), \quad (3.15)$$

and

$$\frac{d\dot{x}}{d\tau} = \frac{b-a}{2}f(x(\tau), \tau) \quad (3.16)$$

Defining two constant vectors

$$X_0 = [2x_0, 0, 0, \dots, 0]^T, X_0 \in R^{N+1} \quad (3.17)$$

$$V_0 = [2\dot{x}_0, 0, 0, \dots, 0]^T, V_0 \in R^{N+1} \quad (3.18)$$

leads to the following steps to obtain the formulations for the position-only MCPI

methods

$$\vec{x}^{k+1} = C_x \vec{\alpha}_x^k \quad (3.19)$$

$$= C_x (C_\alpha \vec{x}^k + X_0) \quad (3.20)$$

$$= C_x \left(C_\alpha \frac{b-a}{2} \vec{x}^k + X_0 \right) \quad (3.21)$$

$$= C_x \left(C_\alpha \frac{b-a}{2} C_x \vec{\alpha}_{vx}^k + X_0 \right) \quad (3.22)$$

$$= C_x \left(C_\alpha \frac{b-a}{2} C_x \left(C_\alpha \frac{b-a}{2} \vec{f}^k + V_0 \right) + X_0 \right) \quad (3.23)$$

$$= C_x C_\alpha \frac{b-a}{2} C_x C_\alpha \frac{b-a}{2} \vec{f}^k + C_x C_\alpha \frac{b-a}{2} C_x V_0 + C_x X_0 \quad (3.24)$$

$$= C_1 \vec{f}^k + C_2 \quad (3.25)$$

with the definitions

$$C_1 = \left(\frac{b-a}{2} C_x C_\alpha \right)^2 \quad (3.26)$$

$$C_2 = \frac{b-a}{2} C_x C_\alpha C_x V_0 + C_x X_0 \quad (3.27)$$

\vec{x} , $\vec{\dot{x}}$, $\vec{\ddot{x}}$, and \vec{f} are vector forms of x , \dot{x} , \ddot{x} , and f evaluated at the $(N+1)$ CGL nodes respectively, and $\vec{\alpha}_x$ and $\vec{\alpha}_{vx}$ are vector forms of the $(N+1)$ Chebyshev coefficients of x and \dot{x} . Comparing Eq. 3.8 with Eq. 3.26, we notice that for linear cases, the convergence condition for the maximum eigenvalues of C_1 and K_1 within the unit circle remains the same when the second order formulations are used.

E. Perturbation Motion MCPI Methods

For a generic dynamic system

$$\frac{d\mathbf{x}(t)}{dt} = \mathbf{f}(t, \mathbf{x}(t)), t \in [a, b] \quad (3.28)$$

with a reference motion described by

$$\frac{d\mathbf{x}_r(t)}{dt} = \mathbf{g}(t, \mathbf{x}_r(t)), t \in [a, b] \quad (3.29)$$

a perturbation motion is defined as

$$\delta\mathbf{x}(t) = \mathbf{x}(t) - \mathbf{x}_r(t) \quad (3.30)$$

leading to the dynamic equation for the perturbation motion as

$$\frac{d\delta\mathbf{x}(t)}{dt} = \mathbf{f}(t, \mathbf{x}_r(t) + \delta\mathbf{x}(t)) - \mathbf{g}(t, \mathbf{x}_r(t)), t \in [a, b] \quad (3.31)$$

One way to improve MCPI methods is to only integrate this perturbation motion $\delta\mathbf{x}(t)$, which is the deviation of the true motion from the reference motion. It is important to note that Eq. 3.31 is an *exact* equation for the perturbation motion. Solving IVPs by this approach can be computationally attractive for several reasons. First, if the reference motion already satisfies the boundary conditions, MCPI methods solve for zero initial and final boundary condition problems, which will simplify the computation. Second, it is possible that a judicious reference motion leads to a small magnitude perturbation motion that requires a small number of iterations to converge. Third but perhaps the most important is that introducing the reference motion brings additional freedom for MCPI methods, by which we may choose $\mathbf{x}_r(t)$ to affect their convergence properties. Although currently we do not have rigorous proofs for these qualitative advantages, the numerical examples in the next section support the utility of these hypotheses.

F. Numerical Examples

1. Small Perturbation from the Sinusoid Motion

Consider a dynamic equation

$$\frac{dy}{dt} = f(y, t) = \cos(t + \epsilon y), y(0) = y_0 \quad (3.32)$$

where the time range is $0 \leq t \leq H$ and the perturbation parameter range is $0 < \epsilon < 1$. Notice for the extreme case that ϵ equals zero, the solution is a simple sinusoidal motion. Fukushima has shown that this problem has an analytical solution as follows [45]

$$y(t) = -\gamma t + \frac{2}{\epsilon} \tan^{-1} \frac{\beta + \sigma \cos \phi}{1 + \sigma \sin \phi} \quad (3.33)$$

where

$$\sigma = \alpha(\sin \phi + \beta \cos \phi) \quad (3.34)$$

$$\phi = \frac{1}{2}(1 - \gamma\epsilon)t \quad (3.35)$$

$$\alpha = \frac{2\epsilon}{1 - \epsilon + \sqrt{1 - \epsilon^2}} \quad (3.36)$$

$$\beta = \frac{1}{1 + \alpha} \tan \frac{\epsilon y_0}{2} \quad (3.37)$$

$$\gamma = \frac{\epsilon}{1 + \sqrt{1 - \epsilon^2}} \quad (3.38)$$

We compare the results by using MCPI methods implemented in MATLAB and by using ODE45, which is a Runge-Kutta 4 – 5 method implemented in MATLAB. ODE45 is simply a convenient (and familiar!) reference to provide a qualitative assessment of the MCPI approach. Three large time intervals (64π , 128π , and 256π) are considered to compare the performance of the two methods with respect to the integration interval length. ϵ is chosen as 0.01 and 0.001 to compare the performance of the two algorithms with respect to the perturbation parameter. The order of

Chebyshev polynomials is adjusted such that the accuracy of the MCPI method is close to or better than ODE45. Figures 8 through 13 show the error distributions of the ODE45 and the MCPI method. Figures 14 through 17 show the polynomial orders and the achieved speedup for two different perturbation parameters.

For qualitative purpose, we note that expanding Eq. 3.32 in ϵ leads to $\frac{dy}{dt} = \cos(t) - \epsilon \sin(t)y + \dots$, so the linear (in y) coefficient is bounded by $\pm\epsilon$. Even though it is not rigorous, we can estimate from the analysis of the constant coefficient linear system in Eq. 3.32 that convergence might be expected if $H < \frac{2}{\epsilon \max(\lambda(C_x C_\alpha))}$. Thus for $\epsilon = 0.01$ with the chosen polynomial $N > 100$, the convergence condition is approximately that H should be less than $\frac{2}{0.05 \times 0.01} \approx 4000$; and for $\epsilon = 0.001$ cases, the convergence domain for H is about 40000. While these estimations may be too optimistic, we verified excellent convergence was actually achieved if $H \leq 320\pi = 1005.3$.

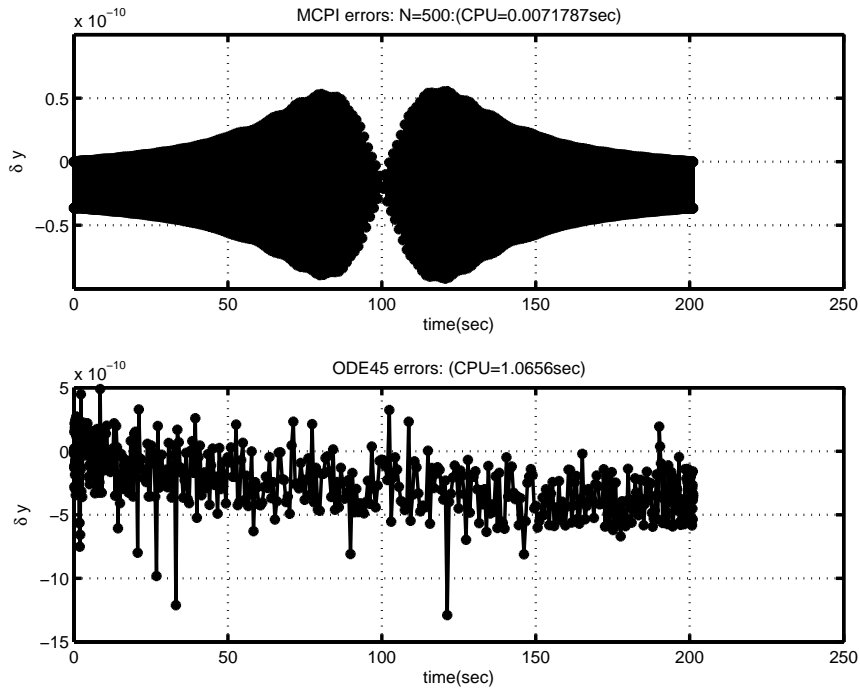


Fig. 8. Integration errors and CPU time comparison ($\epsilon = 0.01$, $H = 64\pi$)

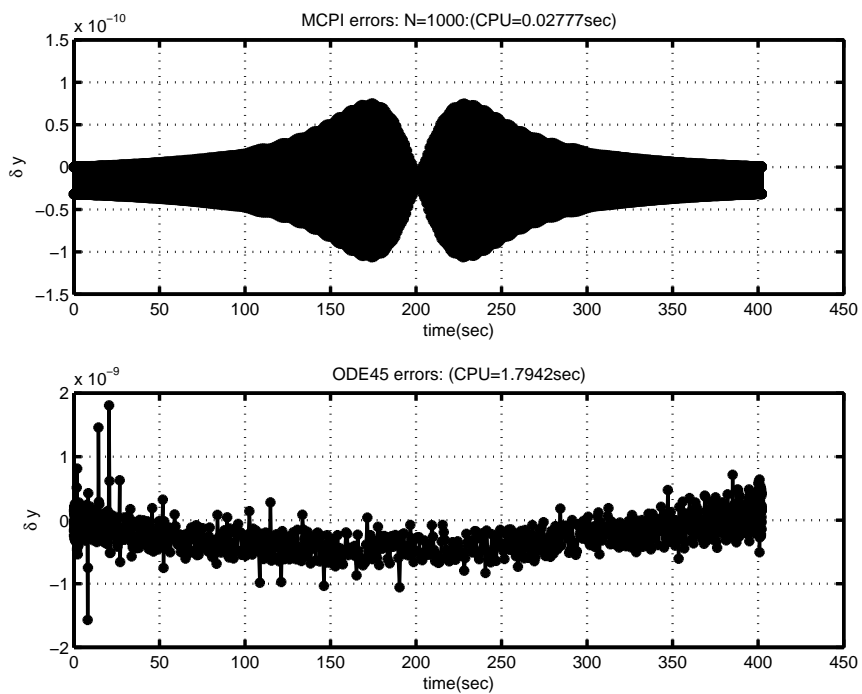


Fig. 9. Integration errors and CPU time comparison ($\epsilon = 0.01$, $H = 128\pi$)

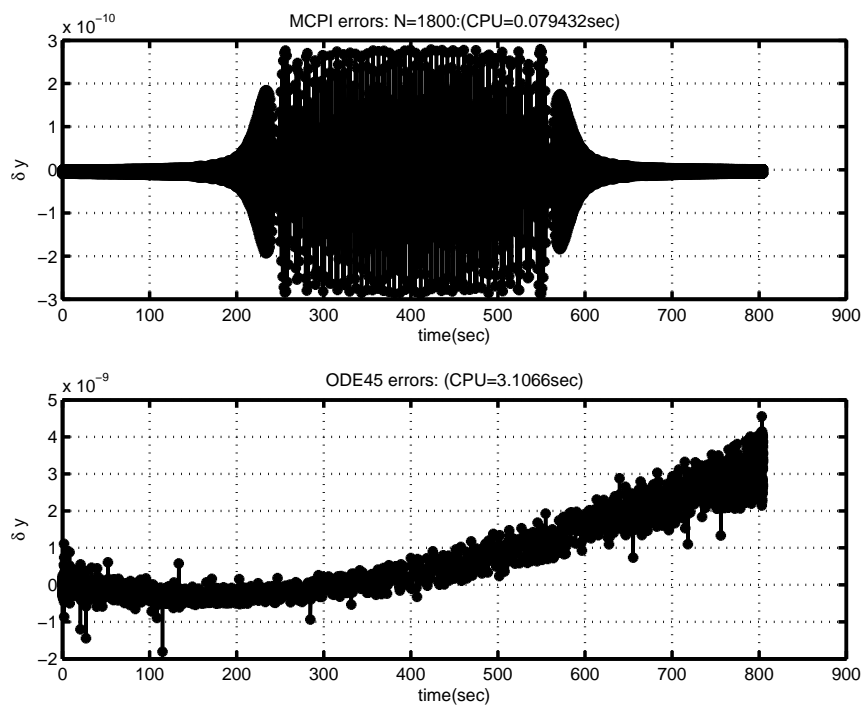


Fig. 10. Integration errors and CPU time comparison ($\epsilon = 0.01$, $H = 256\pi$)

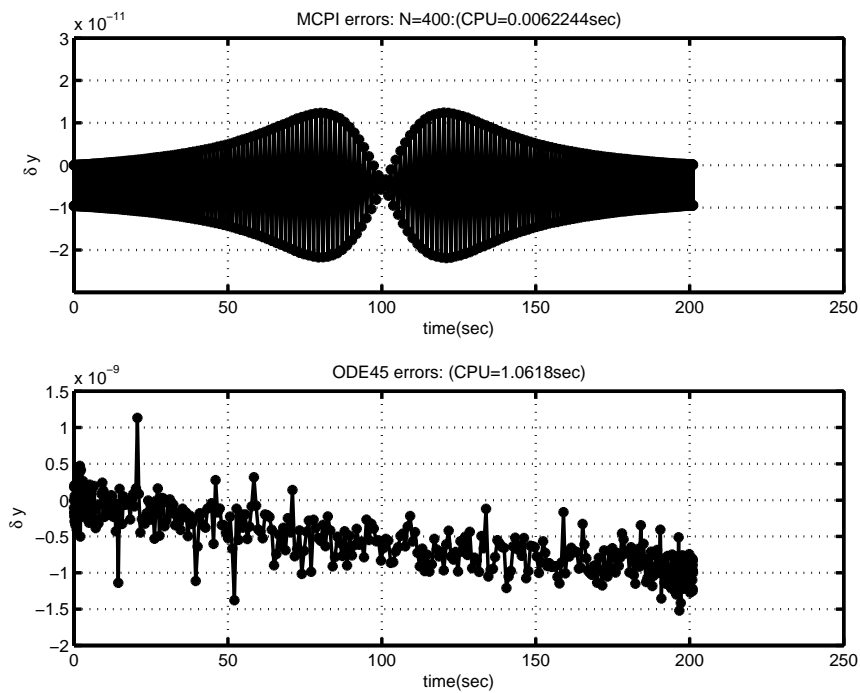


Fig. 11. Integration errors and CPU time comparison ($\epsilon = 0.001$, $H = 64\pi$)

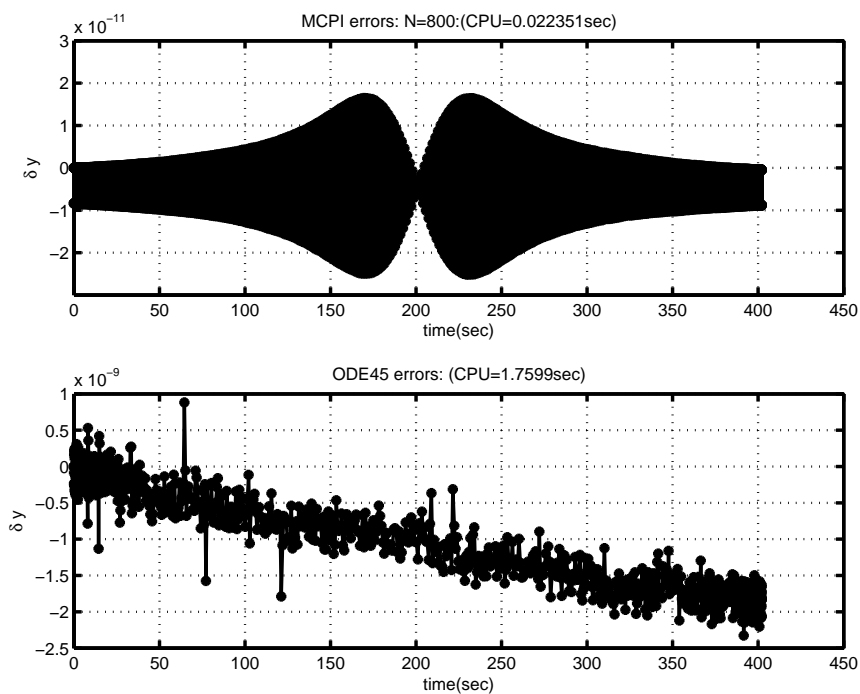


Fig. 12. Integration errors and CPU time comparison ($\epsilon = 0.001$, $H = 128\pi$)

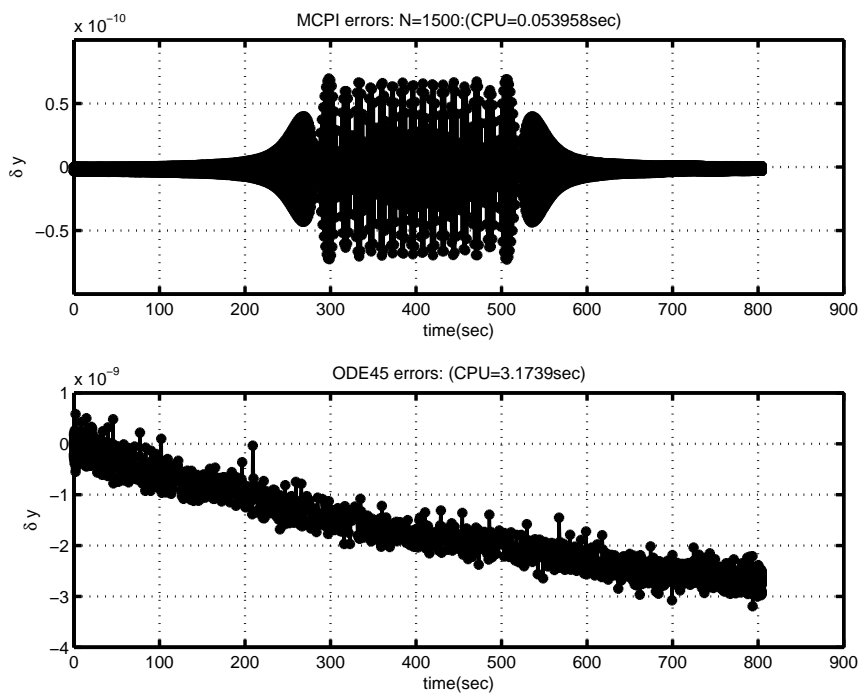


Fig. 13. Integration errors and CPU time comparison ($\epsilon = 0.001$, $H = 256\pi$)

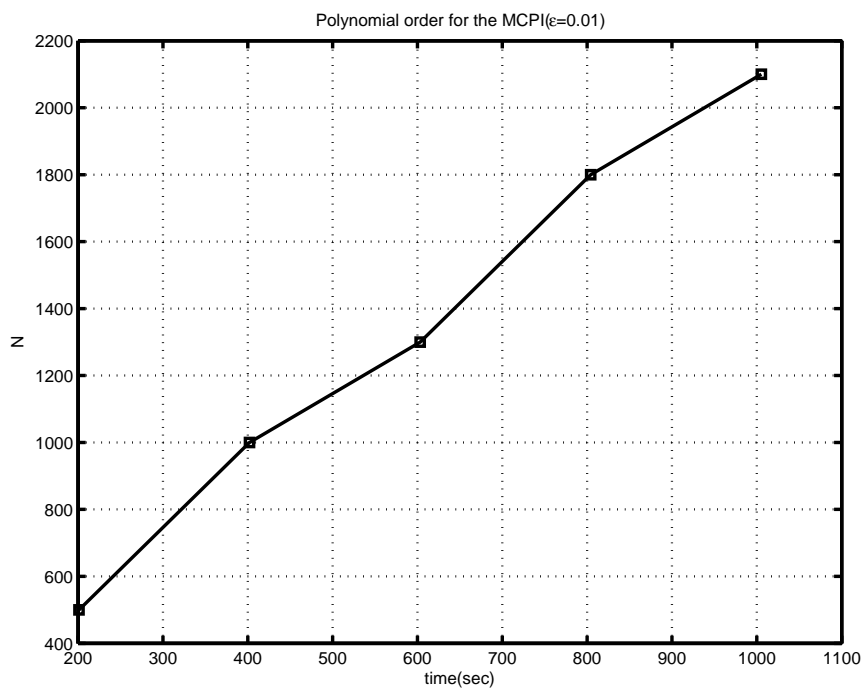


Fig. 14. Polynomial order for the MCPI method ($\epsilon = 0.01$)

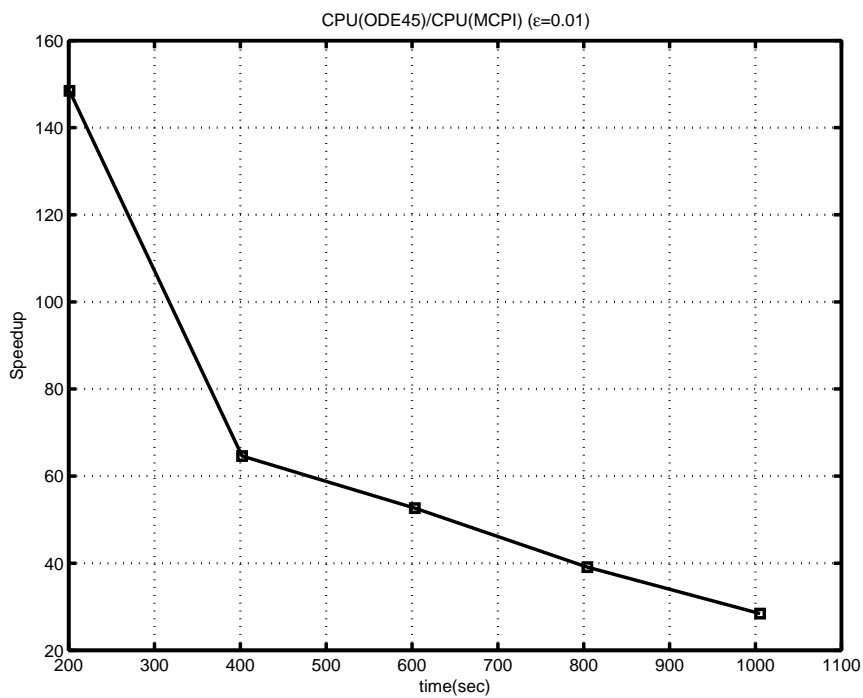


Fig. 15. Speedup by using the MCPI method($\epsilon = 0.01$)

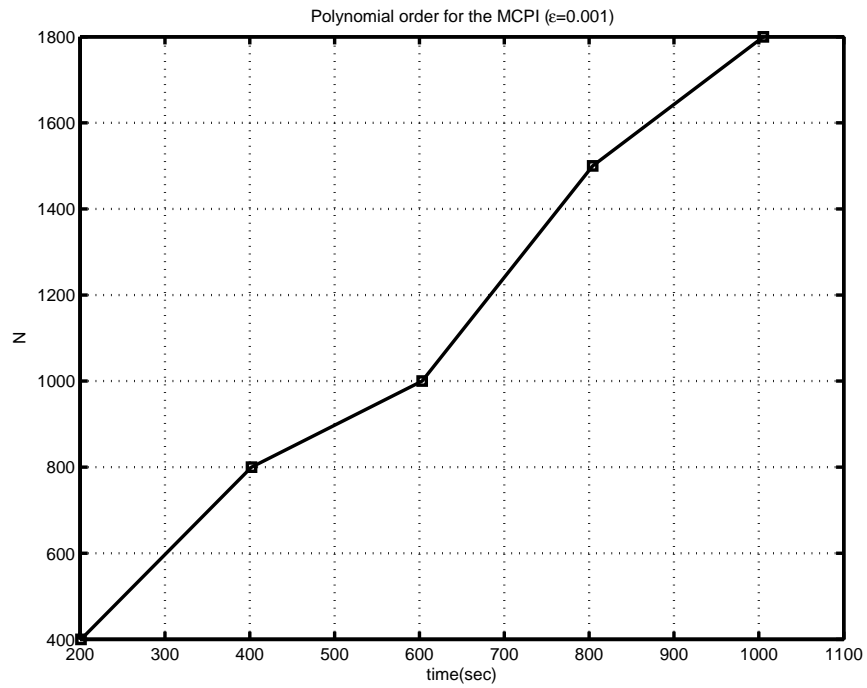


Fig. 16. Polynomial order for the MCPI method ($\epsilon = 0.001$)

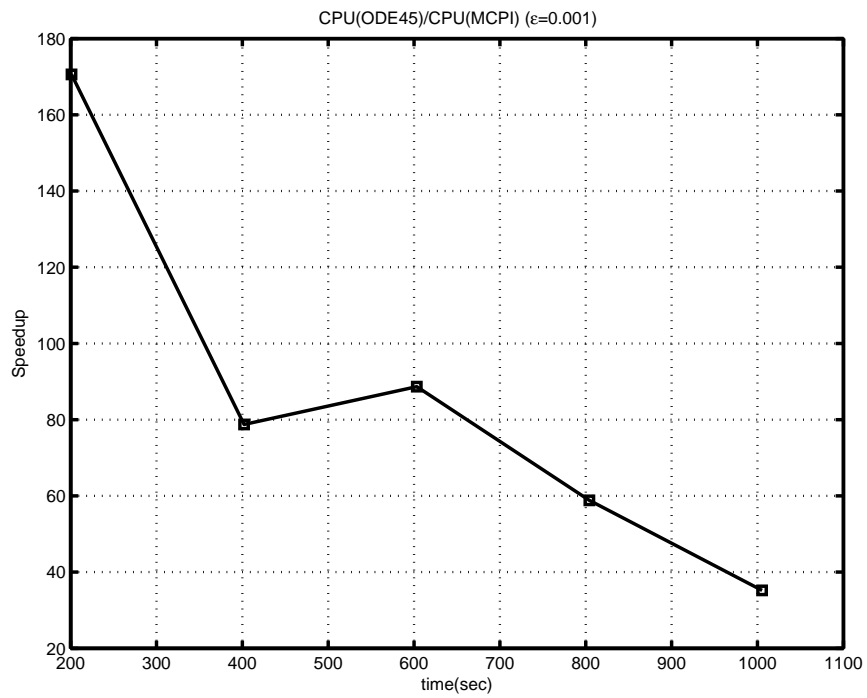


Fig. 17. Speedup by using the MCPI method($\epsilon = 0.001$)

The following conclusions can be drawn

- As both algorithms obtained comparable high accuracy, the CPU time using ODE45 is about 30 to 170 times of the CPU time using MCPI methods.
- While the errors from the reference ODE45 solution have a secular increase, which is a typical result for all forward integration methods, the errors using MCPI method have the maximum values near the middle of the interval and the smallest errors at the boundaries. To graphical precision on a log scale, there is negligible secular error growth, in this example.
- For MCPI methods, the longer the integration interval, the higher order polynomials are required to maintain the accuracy. Also, we note convergence can be obtained up to some problem dependent maximum final time. For linear problem, this maximum can be determined. For nonlinear problems, such as this one, approximation is required. The interval for a practical convergence is significantly greater than 254π (about 128 oscillation periods). Note these solutions do not take advantages of the fact that the long intervals can be subdivided, which will reduce the order of the required polynomial approximation and enable solutions over arbitrary time intervals. We further note that orders of several hundred are feasible, owing to the orthogonality properties and efficient recursions.
- For shorter integration interval, higher speedup is obtained by using MCPI methods, because the orthogonal polynomials converge more efficiently.
- Lower values of ϵ allow lower order Chebyshev polynomials to accurately approximate the solution, essentially because the motion simplifies to sinusoid motion.

2. Satellite Motion Integration by a Piecewise MCPI Approach

Consider a near Earth satellite motion integration problem where only the gravitational force from the Earth is considered. The three dimensional dynamical equations are

$$\ddot{x} = -\frac{\mu}{r^3}x \quad (3.39)$$

$$\ddot{y} = -\frac{\mu}{r^3}y \quad (3.40)$$

$$\ddot{z} = -\frac{\mu}{r^3}z \quad (3.41)$$

where x , y , and z are the three coordinates in some Earth-centered inertial reference frame; r is the distance of the satellite from the Earth; μ is the Earth gravitational constant and is chosen as $3.986 \times 10^5 km^3/s^2$. Even though these equations are non-linear, they can be solved analytically. Here we solve the two-body Keplerian motion using the F and G approach [42] to provide an exact analytical baseline for verifying the results. The initial position and velocity are

$$\mathbf{r}(t_0) = [-464.856, 6667.880, 574.231]^T km \quad (3.42)$$

$$\mathbf{v}(t_0) = [-2.8381186, -0.7871898, 7.0830275]^T km/s \quad (3.43)$$

which lead to a near circular orbit with the classical orbital elements

$$a = 6644.754 km \quad (3.44)$$

$$e = 0.0099865 \quad (3.45)$$

$$i = 1.1866 rad \quad (3.46)$$

$$\Omega = 1.6057 rad \quad (3.47)$$

$$\omega = -2.7805 rad \quad (3.48)$$

$$T_p = 5.3905 \times 10^3 sec \quad (3.49)$$

where a is the semimajor axis, e is the eccentricity, i is the inclination angle, Ω is the longitude of the ascending node, ω is the argument of periapsis, and T_p is the orbit period. We compare the performance of the piecewise MCPI method with a reference ODE45 solution and the results shown here are obtained by using one complete orbit period as the segment step size for MCPI methods. Figure 18 compares the history of the errors for ten orbit revolution integration. The MCPI solution is shown to have up to three orders of magnitude better accuracy than the reference ODE45 solution. The CPU time using the two methods is shown in Fig. 19. Figure 20 demonstrates that the speedup of the MCPI method over the reference ODE45 solution is over one order of magnitude. Achieving both better accuracy and significant speedup, we can conclude that MCPI methods are well-suited with orbit propagation problems. While these present results are for the two-body problem (to allow rigorous error discussion, since we have an analytical solution) as will be evident in Section 4 below, inclusion of dominant gravity do not alter the conclusion that substantial speedups are achieved, while maintaining high accuracy.

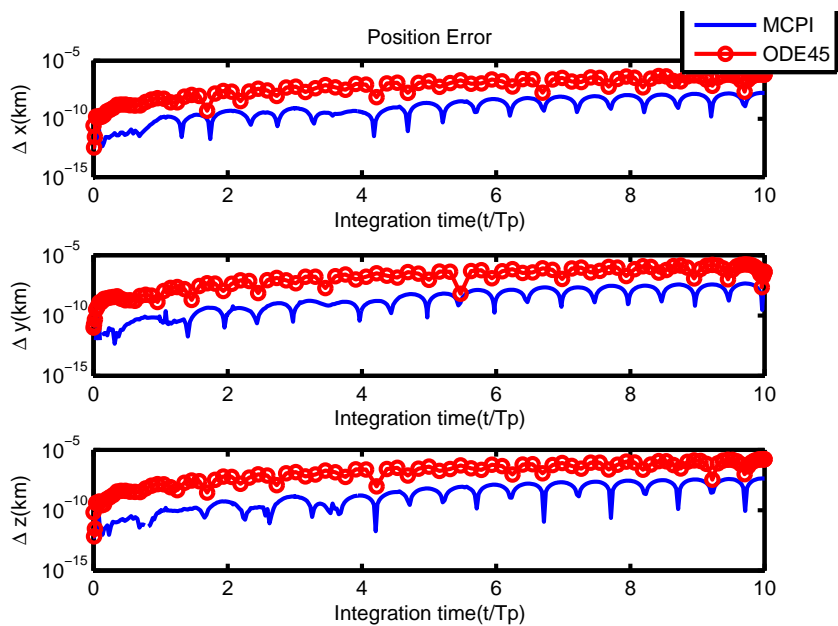


Fig. 18. Error history of MCPI and ODE45 for ten orbit (two-body problem)

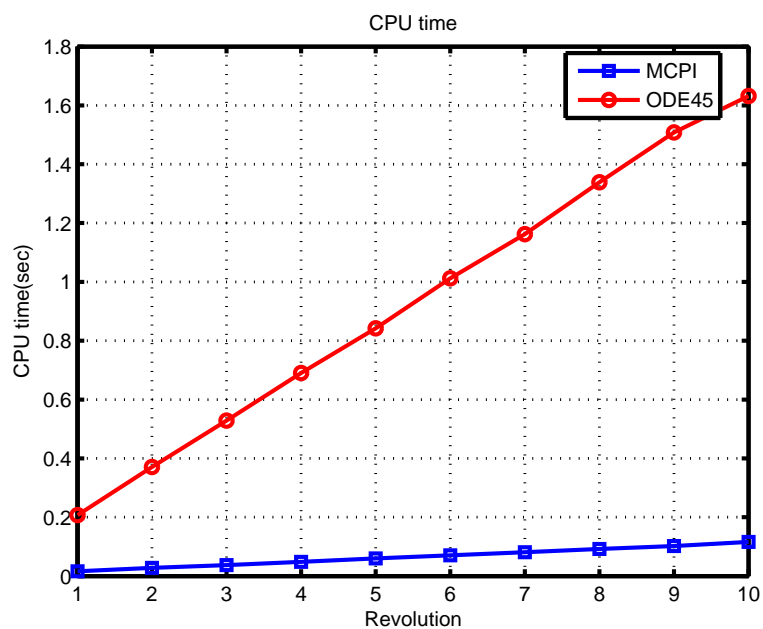


Fig. 19. CPU time of MCPI and ODE45 (two-body problem)

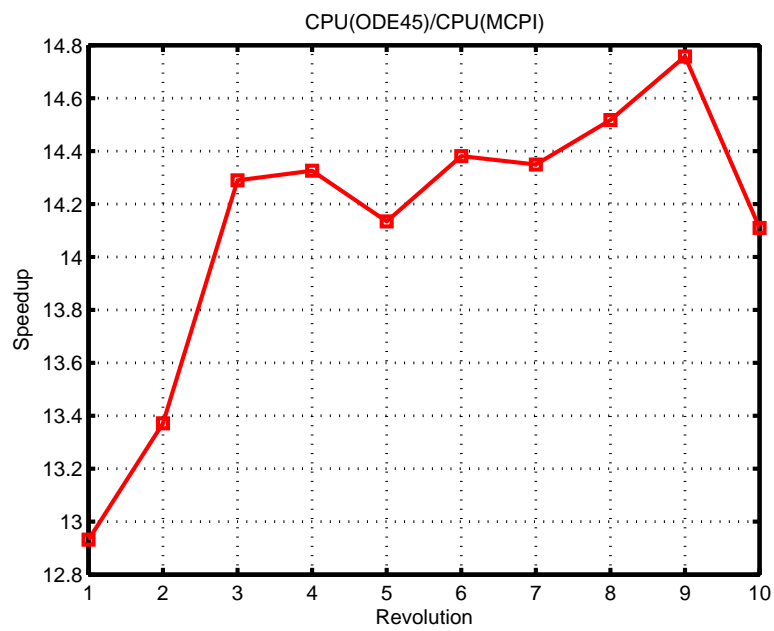


Fig. 20. Speedup of MCPI over ODE45 (two-body problem)

3. Results of the Position Only MCPI Method

First, we apply the position only MCPI method to the satellite motion integration problem discussed in Section 2. The satellite motion is integrated for 16 orbits, leading to an integration time about one day. Figure 21 shows the CPU time for the reference ODE45 solution, the MCPI method, and the position only MCPI method. The speedup of the position only MCPI method over the reference ODE45 solution is shown in Fig. 22, and its speedup over the original MCPI method is shown in Fig. 23. These plots show that the longer time integration, the more speedup is obtained by using the position only MCPI method.

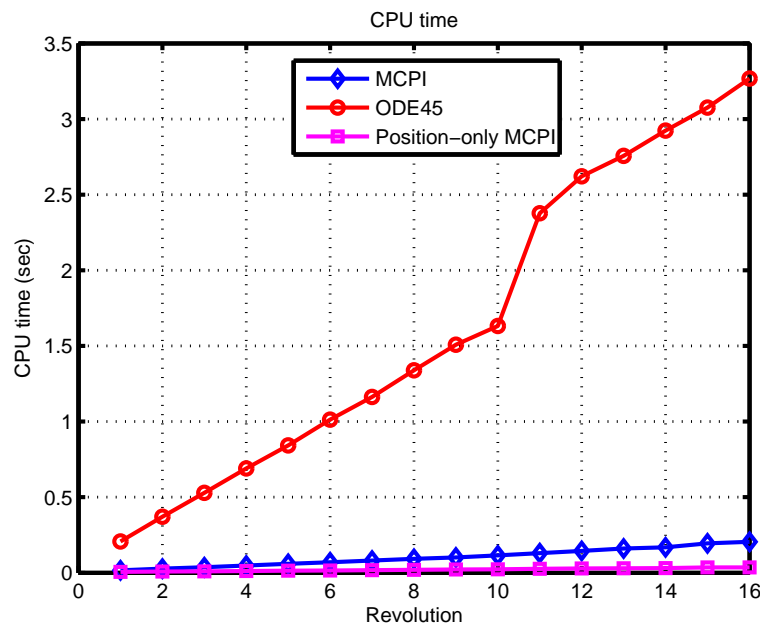


Fig. 21. CPU time of ODE45, MCPI, and position only MCPI (two-body problem)

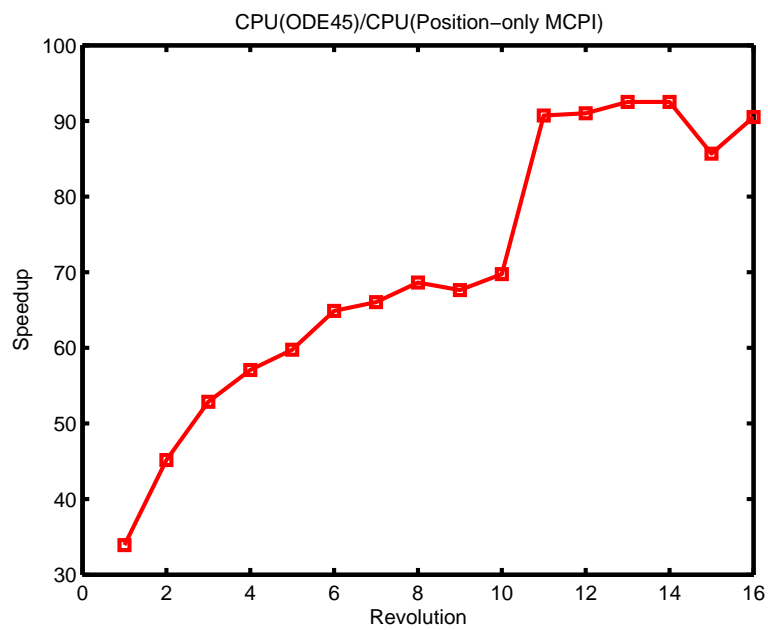


Fig. 22. Speedup of position only MCPI over ODE45 (two-body problem)

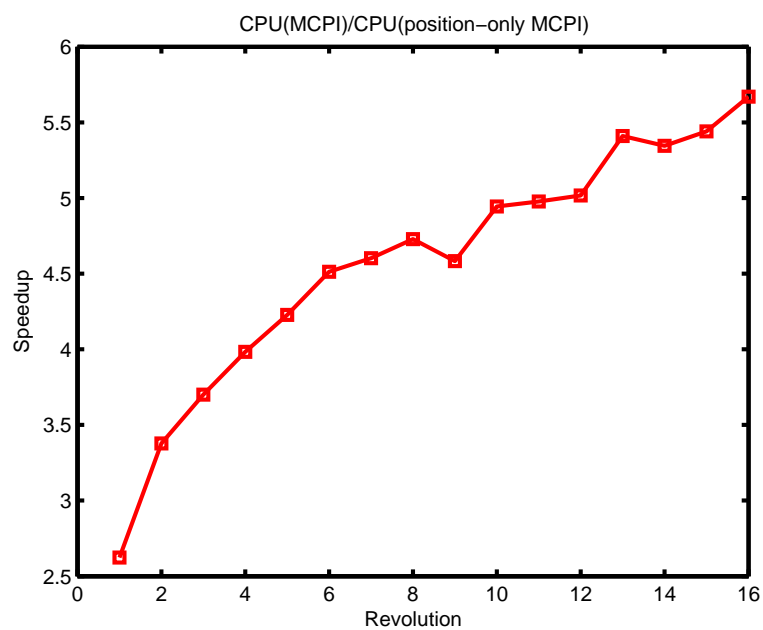


Fig. 23. Speedup of position only MCPI over the original MCPI (two-body problem)

Next, we use the position only MCPI method to solve a planar case three-body problem. We study this problem to show the applicability of MCPI methods for highly coupled and higher dimensional problems. The dynamic equations are

$$\ddot{x}_1 = -\frac{\mu_2}{r_{12}}x_{12} - \frac{\mu_3}{r_{13}}y_{13} \quad (3.50)$$

$$\ddot{y}_1 = -\frac{\mu_2}{r_{12}}y_{12} - \frac{\mu_3}{r_{13}}y_{13} \quad (3.51)$$

$$\ddot{x}_2 = \frac{\mu_1}{r_{12}}x_{12} - \frac{\mu_3}{r_{23}}y_{23} \quad (3.52)$$

$$\ddot{y}_2 = \frac{\mu_1}{r_{12}}y_{12} - \frac{\mu_3}{r_{23}}y_{23} \quad (3.53)$$

$$\ddot{x}_3 = \frac{\mu_1}{r_{13}}x_{13} + \frac{\mu_2}{r_{23}}y_{23} \quad (3.54)$$

$$\ddot{y}_3 = \frac{\mu_1}{r_{13}}y_{13} + \frac{\mu_2}{r_{23}}y_{23} \quad (3.55)$$

where μ_1 , μ_2 , and μ_3 are the gravitational constants of the three bodies; (x_1, y_1) , (x_2, y_2) , and (x_3, y_3) are the three inertial coordinates for the three bodies, and

$$x_{12} = x_1 - x_2 \quad (3.56)$$

$$x_{13} = x_1 - x_3 \quad (3.57)$$

$$x_{23} = x_2 - x_3 \quad (3.58)$$

$$y_{12} = y_1 - y_2 \quad (3.59)$$

$$y_{13} = y_1 - y_3 \quad (3.60)$$

$$y_{23} = y_2 - y_3 \quad (3.61)$$

The three body masses are $m_1 = 5.967 \times 10^{23}$ kg, $m_2 = 7.53 \times 10^{22}$ kg, and $m_3 =$

5.967×10^{23} kg. The initial position are

$$x_1(t_0) = -1.2996 \times 10^8 km \quad (3.62)$$

$$y_1(t_0) = -0.4502 \times 10^8 km \quad (3.63)$$

$$x_2(t_0) = 8.7004 \times 10^8 km \quad (3.64)$$

$$y_2(t_0) = -0.4502 \times 10^8 km \quad (3.65)$$

$$x_3(t_0) = 3.7004 \times 10^8 km \quad (3.66)$$

$$y_3(t_0) = 8.2101 \times 10^8 km \quad (3.67)$$

and the initial velocities are

$$\dot{x}_1(t_0) = -15.3344 km/s \quad (3.68)$$

$$\dot{y}_1(t_0) = -25.6286 km/s \quad (3.69)$$

$$\dot{x}_2(t_0) = 151.0113 km/s \quad (3.70)$$

$$\dot{y}_2(t_0) = 113.9520 km/s \quad (3.71)$$

$$\dot{x}_3(t_0) = -53.0418 km/s \quad (3.72)$$

$$\dot{y}_3(t_0) = 188.2210 km/s \quad (3.73)$$

These parameters and initial conditions are calculated from the example used by Schaub and Junkins [42]. For these special initial conditions, following Lagrange, the motion is exactly solvable and an analytical equilateral solution using the F and G approach provides a baseline for the verification. More general initial conditions do not have an analytical solution, but energy and angular momentum are conserved in general. The obtained triangular configuration is highlighted in Fig. 24 and the invariant shape of the equilateral triangle is obvious. The longest integration time we use is $3 \times 10^6 sec \approx 34.72 days$. For this case, CPU time using a reference ODE45

solution is 0.09sec while for position-only MCPI method, CPU time is 0.007sec by using polynomial order of one hundred. The relative position errors of the position only MCPI method and the reference ODE45 solution are shown in Fig. 25. The capability to have large “step size” for MCPI methods is further demonstrated here. The speedup of the position only MCPI method over the reference ODE45 solution with respect to different integration times is shown in Fig. 26.

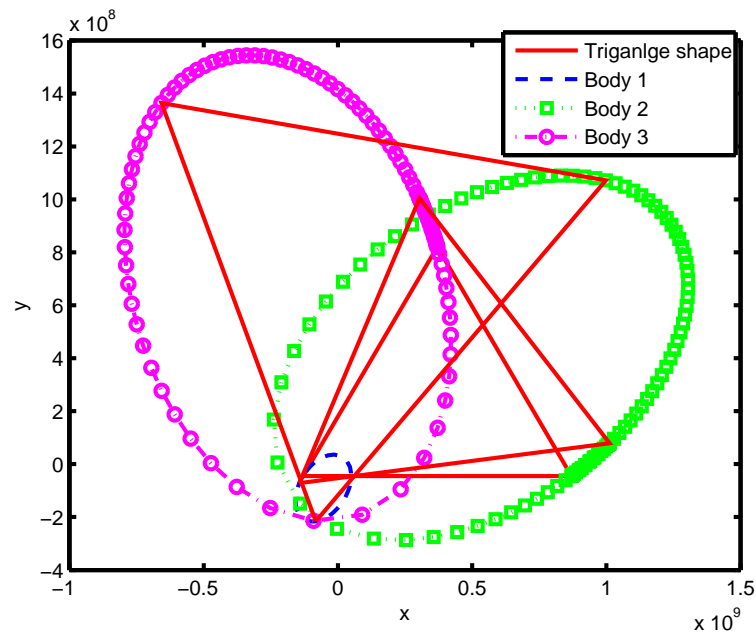


Fig. 24. Configuration of the three-body motion

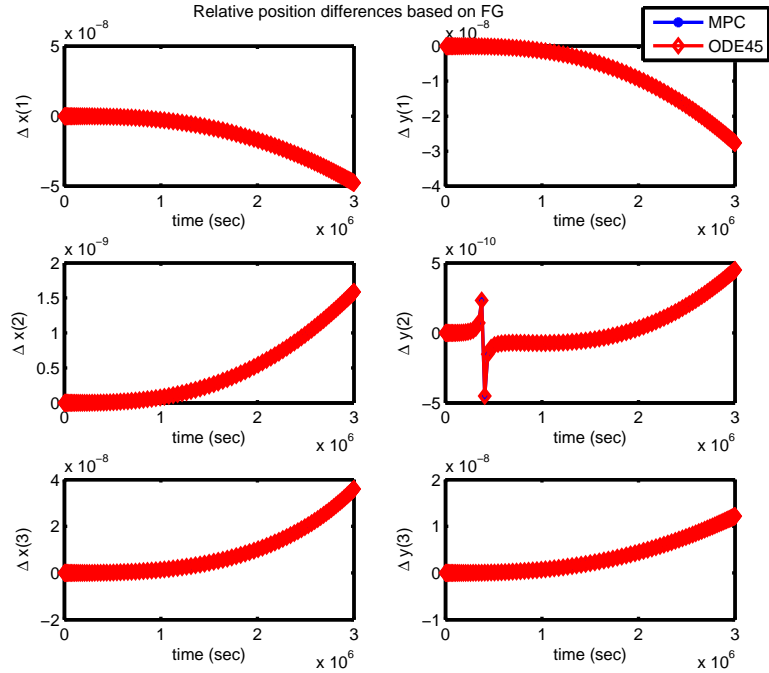


Fig. 25. Relative position errors based on FG solution (three-body motion)

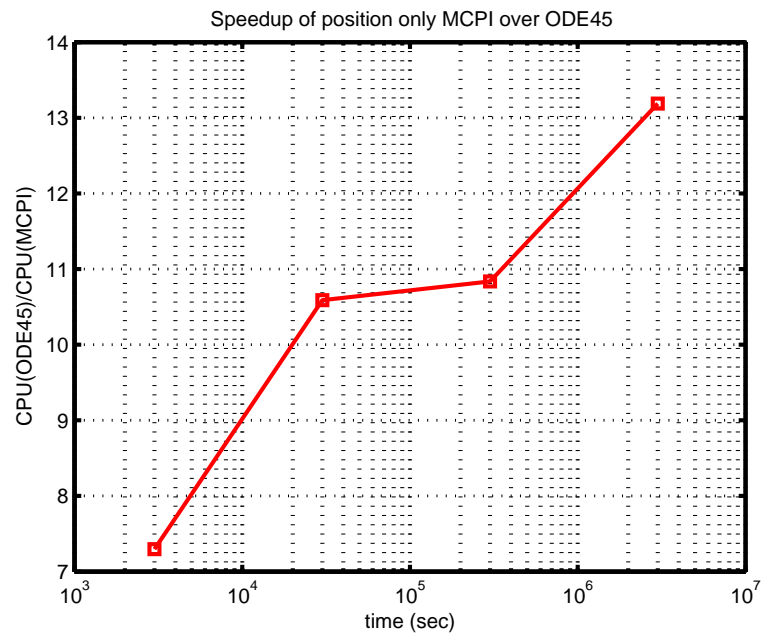


Fig. 26. Speedup of position only MCPI over ODE45 (three-body motion)

4. Zonal Harmonic Perturbation Satellite Motion Propagation Problems

In the celestial mechanics field, very often the orbit calculations require inclusion of many complicated perturbation terms to model perturbation accelerations in addition to the dominant force. For perturbed two-body problems, the relative motion between the two body is described by

$$\ddot{\mathbf{r}} = -\frac{\mu}{r^3}\mathbf{r} + \mathbf{a}_d \quad (3.74)$$

where \mathbf{r} is the relative position between the two bodies and \mathbf{a}_d is the disturbance acceleration from the perturbation forces. Equation 3.74 can be integrated directly by conventional ODE solvers, which is the often called Cowell's method. Another way is to only integrate the perturbation motion while using the unperturbed two-body solution as the reference motion, which is usually referred to as Encke's method.

Here we use the satellite motion integration problem presented in Section 2 as an example to show the benefit of using the perturbation only MCPI method. The dominant force is the inverse-square gravity force and we include the zonal harmonic perturbation forces up to the fifth order [42]. Including zonal harmonic perturbations up to the order of k leads to the dynamic equation for the perturbation motion as

$$\delta\ddot{\mathbf{r}} = -\frac{\mu}{r^3}\mathbf{r} + \sum_{i=2}^{i=k}\mathbf{a}_d^i + \frac{\mu}{r_r^3}\mathbf{r}_r \quad (3.75)$$

where \mathbf{a}_d^i is the i^{th} order perturbation acceleration. Figure 27 illustrates the CPU time using different approaches. In the figure, the method called MCPI with the unperturbed initial guess means that the MCPI method starts with an initial guess which is the solution from the unperturbed two-body problem; the method called position only MCPI means the reference motion is the unperturbed two-body solution and the MCPI method only solves for the perturbation motion. Figure 28 shows the speedup of three MCPI methods over the reference ODE45 solution. The perturbation only

MCPI method achieves the most speedup. Also notice as the number of perturbation terms increase, the speedup from using MCPI methods drops. At first glance, this figure suggests that the advantages are lost as the complexity of the acceleration model increases. However, as will be evident in the parallel computation discussion later, the opposite is true. This is because, along any iterative trajectory the acceleration is an explicit function of time and therefore each term of the model can be computed in parallel threads. The performance of MCPI methods can be significantly improved by two obvious approaches: (i) using parallel computation; and (ii) compiling the current code before running instead of being interpreted by MATLAB during execution. The first of these approaches is conclusively verified in Chapter VI.

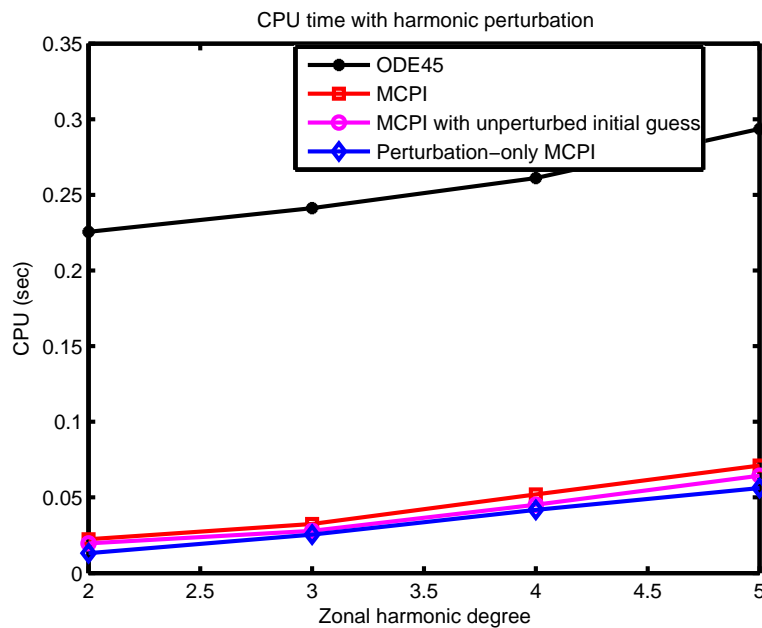


Fig. 27. CPU time of different MCPI and ODE45 (two-body motion)

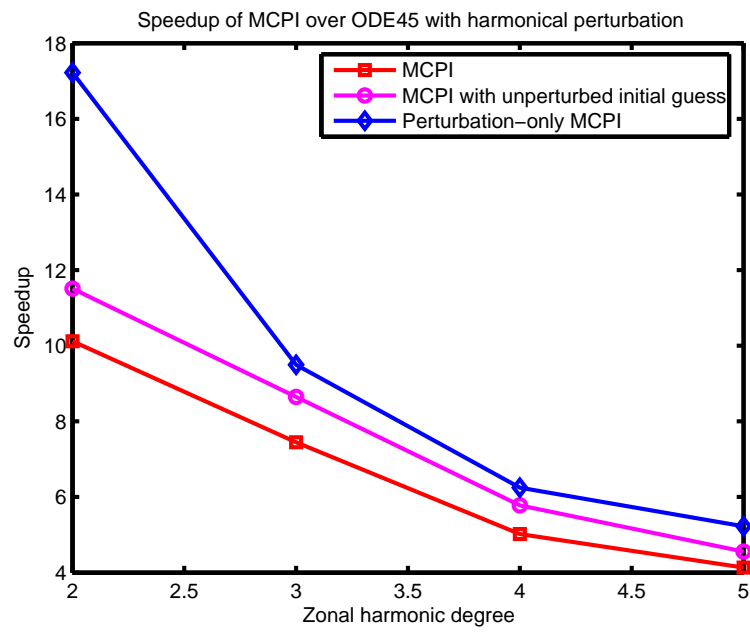


Fig. 28. Speedup of different MCPI and ODE45 (two-body motion)

G. Parameters Affecting the Performance of MCPI Methods

Computational time and accuracy are the main criteria we choose to evaluate the performance of MCPI methods. The parameters that may affect the performance of MCPI methods are: (i) the order of the approximate Chebyshev polynomials; (ii) the stopping criterion; and (iii) the finite segment step size which needs to be decided when the piecewise approach is used for long time integration. The general rule is that higher order polynomials and smaller tolerance to end the iteration will lead to better accuracy solutions but require longer computation time. Given that the segment step size is smaller than the maximum that leads to convergence of the process shown in Fig. 3, an additional crucial tradeoff must be resolved. The question is: What is the optimal polynomial order that will be as efficient as possible subject to the constraint that required accuracy is achieved? Clenshaw and Norton pointed out this optimal order will generally be difficult to know in advance, and suggested starting with lower order and increasing the polynomial order when necessary [30]. This suggests an adaptive strategy qualitatively analogous to the automatic step size control in forward integration methods for solving differential equations. Steele developed a formulation that needs to be iteratively solved to obtain the valid interval to approximate the solutions [46]. However, the proposed formulation in its current stage of evolution remains difficult to apply to vector cases. From our literature review, the most popular approach is to find the optimal order through experiments [32, 45]. Furthermore, there have been no historical discussions on how to choose the segment step size. However, alone we found important insight which are near-conclusive for the case of linear systems and provide important insight for nonlinear systems. Note the segment step size has a double-edged sword effect on the performance of the piecewise MCPI method, and obviously couples strongly with choosing the order of approximation

and the resulting accuracy of the approximation. Through our experiments, trial and error remains the most efficient approach to date, with some appropriate insight obtained from linearizing the system. However, we anticipate adaptive algorithms for adjusting the order and segment length will emerge from future studies.

Here we use the two-body problem discussed in Section 2 to provide some insight on how to choose these parameters. The segment step size is defined as the finite interval length over which we use one set of Chebyshev polynomials to approximate the solution, and is denoted as h in the figures. The stopping criterion for the MCPI method is chosen as $\epsilon = 10^{-5}$, which means the iteration stops whenever both the corrections at the previous iteration and the corrections on the stopping iteration are less than 10^{-5} . In this problem, this tolerance is found to guarantee the maximum errors along all the three directions are less than one meter.

Figure 29 shows the CPU time versus polynomial orders as the segment step sizes are chosen for three different values. It shows that once the stopping criterion and segment step size are fixed, as the order of the polynomial increases, the CPU time goes up as well. Additionally, it appears that there exists an optimal segment step size that achieves the minimum computation time. For the three segment step sizes we compare, the three hour segment step size is the best time segment length to obtain the most speedup. The reason is: as the segment length gets longer, less number of segments is required (the overall solutions are obtained by patching the individual solutions together), which could reduce the computation time. However, as the segment length gets much longer, higher order polynomials are required to approximate the solutions, which may also require more and more expensive iterations to satisfy the stopping criterion and thus increase the computation time. An optimal segment length balances both the number of segments and the number of required iterations.

Figure 30 shows the maximum distance errors with respect to the order of polynomials as the segment step size is chosen for three different values. The F and G solutions are calculated beforehand to provide a baseline. First, we observe that for a fixed segment step size, higher order polynomials usually bring better accuracy. Second, we notice in most cases, the larger the segment step size, the better accuracy can be achieved, which is because the accumulated errors are reduced by reducing the number of segments. Third, the special cases are that the accuracy drops for $h = 1.5\text{hour}$ and $h = 6\text{hour}$ when we choose the order as 400, which are because the accuracy and computation are coupled issues. In fact, in these cases, higher accuracy can be achieved if we set the stopping criterion higher. However, for a fixed stopping criterion, the iteration can exit earlier before the MCPI method achieves its best accuracy.

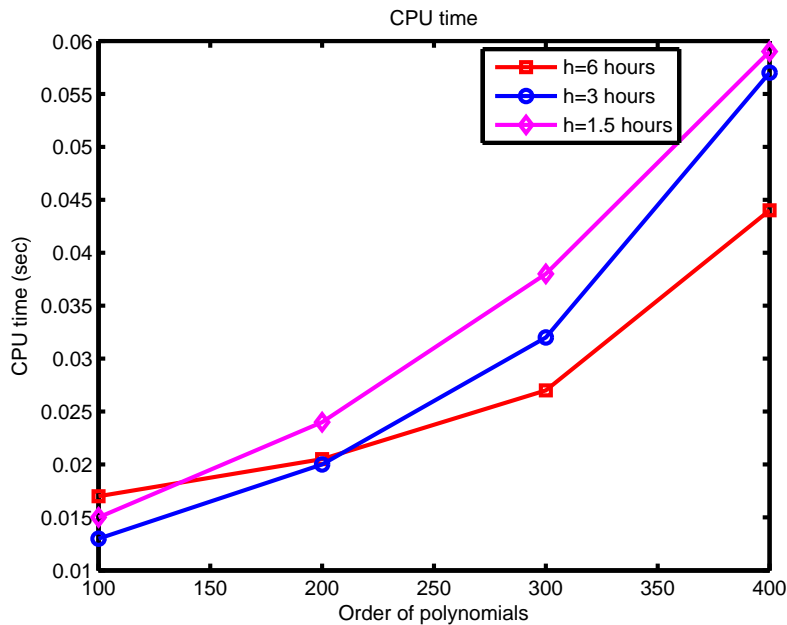


Fig. 29. CPU time comparison (two-body problem)

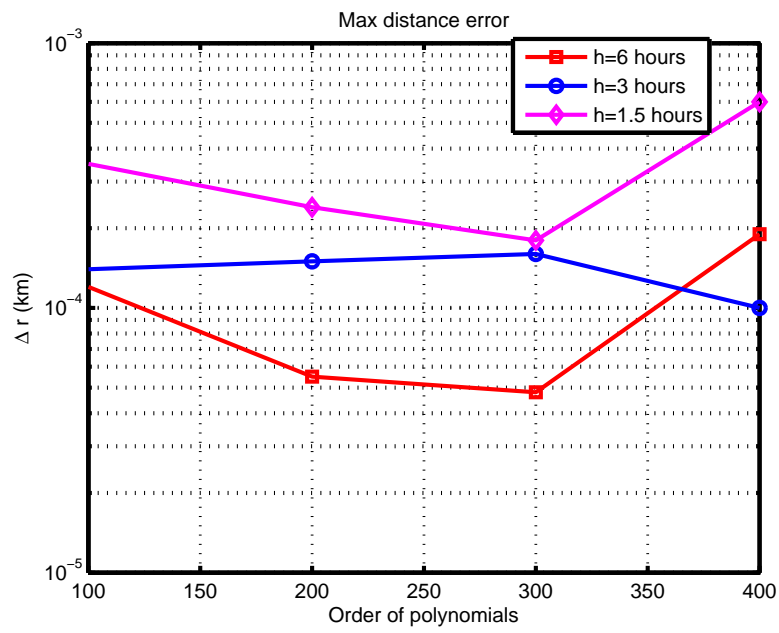


Fig. 30. Max distance error comparison (two-body problem)

From these experimental results, we confirm and quantify the expected results that the computation time and accuracy of MCPI methods are dependent on the chosen polynomial order, segment step size, and stopping criterion. These issues sound analogous to the variable step size, variable order forward integration methods. The obvious differences are that the practical order of MCPI methods can be much higher than forward integration methods and the segment step size of MCPI methods is usually significantly larger.

Although choosing optimal polynomial order and segment step size remains challenging, finding some suboptimal numbers through experiments, which can satisfy the specified requirement, is relatively easy in practice. From our experiments, there are typically large sets of feasible suboptimal decisions in MCPI methods that provide very substantial advantages for this approach.

H. Summary

Although the convergence domain for the original MCPI methods is difficult to know in advance, MCPI methods can efficiently solve many important IVPs over arbitrarily long time intervals by utilizing the piecewise approach introduced in this chapter. Furthermore, the accumulated errors through the piecewise MCPI methods can be significantly smaller than the usual forward integration methods. MCPI methods also provide an “inbuilt” interpolation, since an orthogonal function approximation of the state trajectory is obtained. The proposed MCPI methods are shown to achieve both high accuracy and time efficiency for the orbital propagation problems. The speedup achieved from using the position only MCPI method and perturbation only MCPI method is also shown to be substantial. Regardless of the fact that the global convergence for all problems can not be guaranteed, the evident applicability of MCPI

methods for solving these problems makes a clear statement regarding the importance of the methodology. MCPI methods are shown to obtain speedup over ODE45 reference solutions for all the studied examples, although the magnitude of the speedup is dependent on the problems themselves. We also provide some insight on how to choose the order of Chebyshev polynomials and how to choose the segment step size for the piecewise approach. Next chapter will study using the MCPI methodology to solve BVPs.

CHAPTER IV

MODIFIED CHEBYSHEV-PICARD ITERATION METHODS FOR SOLUTION
OF BOUNDARY VALUE PROBLEMS

A. Introduction

In this chapter, we discuss applications of the MCPI methodology developed in Chapter II to solve two-point boundary value problems (BVPs). First, three forms of BVPs for the second order dynamic systems are introduced due to the historical importance of these problems. Next, a linear second order problem is presented to illustrate the fundamentals of applying MCPI methods for solving BVPs. An MCPI method is then used to solve the classical Lambert's problem, and its performance is compared with several conventional methods. An optimal trajectory design problem is then presented to illustrate practical implications of the procedure and the advantage of using MCPI methods for solving optimal control problems.

B. Three Forms of BVPs for Second Order Systems

Consider the second order ordinary differential equation

$$\ddot{\mathbf{x}}(t) = \mathbf{f}(t, \mathbf{x}, \dot{\mathbf{x}}) \quad (4.1)$$

subject to the requirement that $\mathbf{f}(t, \mathbf{x}, \dot{\mathbf{x}})$ is continuous and satisfies a uniform Lipschitz condition

$$|f(t, \mathbf{x}_1, \dot{\mathbf{x}}_1) - f(t, \mathbf{x}_2, \dot{\mathbf{x}}_2)| \leq L|\mathbf{x}_1 - \mathbf{x}_2| + K|\dot{\mathbf{x}}_2 - \dot{\mathbf{x}}_1| \quad (4.2)$$

where K and L are two finite, non-negative constants. Here $|\cdot|$ represents some distance measurement such as the Euclidean norm of the vector [47]. There exist

three different ways to constrain the boundary conditions:

$$\mathbf{x}(a) = A, \mathbf{x}(b) = B \quad (4.3)$$

or

$$\mathbf{x}(a) = A, \dot{\mathbf{x}}(b) = D \quad (4.4)$$

or

$$\dot{\mathbf{x}}(a) = D, \mathbf{x}(b) = B \quad (4.5)$$

We classify problems in the form of Eqs. 4.1 and 4.3 as two-point BVPs of the first kind, in the form of Eqs. 4.1 and 4.4 as two-point BVPs of the second kind, and in the form of Eqs. 4.1 and 4.5 as two-point BVPs of the third kind. Similar definitions have been used by Craats [47]. Although most of the literatures only consider BVPs of the first kind [8, 9, 10, 11, 12] (discussed in Section 2 of Chapter I), this chapter will show that the proposed MCPI methods are applicable to all three kinds of BVPs.

C. Applying MCPI Methods to a Linear Second Order System

Here, we use a second order linear system as an example to illustrate the methodology of using MCPI methods for solving BVPs. For a differential equation of the form

$$\ddot{y}(\tau) + \lambda^2 y(\tau) = 0, -1 \leq \tau \leq 1 \quad (4.6)$$

we consider four possible cases. In the first case, the constraints for both y and its derivative \dot{y} are given at the final time $\tau = 1$. Although this case can be categorized in initial value problems, we give special attention to it here because it is different from traditional IVPs where the conditions are defined at the initial time. The other

three cases are BVPs of the three kinds discussed in Section B. For all the cases, we choose the parameter $\lambda = 0.5$, and the order of Chebyshev polynomials as 100.

1. Solving Final Value Problems

In this case, $y(1)$ and $\dot{y}(1)$ are known. After transforming the second order differential equation to its first order form, we assume the following Chebyshev polynomial approximations for y and \dot{y}

$$y(\tau) = \sum_{k=0}^{k=N} \alpha_k T_k(\tau) \quad (4.7)$$

$$\dot{y}(\tau) = \sum_{k=0}^{k=N} \beta_k T_k(\tau) \quad (4.8)$$

Similar to the approaches used to solve IVPs, the coefficients α_k and β_k are updated as follows

$$\vec{\alpha} = C_\alpha \dot{\vec{y}} \quad (4.9)$$

$$\vec{\beta} = C_\alpha (-\lambda^2 \vec{y}) \quad (4.10)$$

where $\vec{\alpha}$ and $\vec{\beta}$ are the $(N + 1)$ coefficients α_k and β_k written in the vector form; $\dot{\vec{y}}$ and \vec{y} are the velocity and position evaluated at the CGL points.

The final conditions are

$$y(1) = \sum_{k=0}^{k=N} \alpha_k T_k(1) \quad (4.11)$$

$$y(-1) = \sum_{k=0}^{k=N} \beta_k T_k(1) \quad (4.12)$$

which require the zero order coefficients to be calculated through

$$\alpha_0 = 2y(1) - 2 \sum_{k=1}^{k=N} \alpha_k \quad (4.13)$$

$$\beta_0 = 2\dot{y}(1) - 2 \sum_{k=1}^{k=N} \beta_k \quad (4.14)$$

The iterative procedures for solving for the \vec{y} and $\vec{\dot{y}}$ coefficients are analogous to the way we solve IVPs. The errors of the solutions for this case are shown in Figs. 31 and 32. The boundary conditions at the final time are shown to be satisfied with high accuracy, which is one of the advantages of MCPI methods, although in this case the simplicity of the linear system makes the result unsurprising.

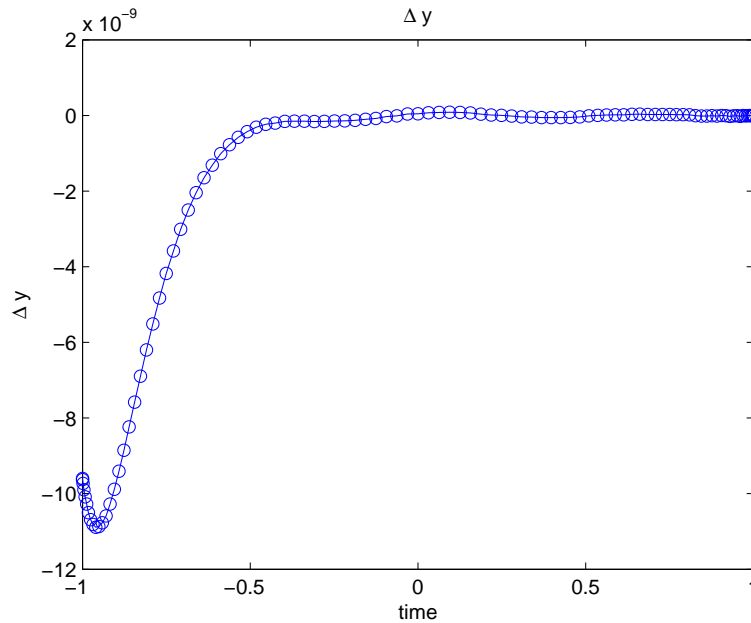


Fig. 31. Integration errors of y (final value problem)

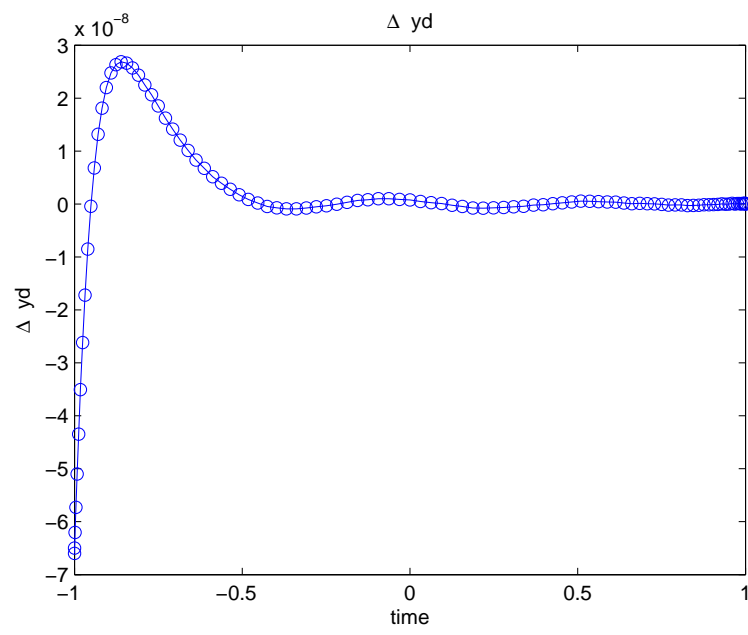


Fig. 32. Integration errors of \dot{y} (final value problem)

2. Solving BVPs of the First Kind

In this case, either both $y(-1)$ and $y(1)$ are known or both $\dot{y}(-1)$ and $\dot{y}(1)$ are known. These boundary conditions are used to constrain both the zero-order and the first-order Chebyshev coefficients for the state elements whose boundaries are constrained.

Consistent with Eq. 2.15 and 2.16, for a general state x , where x can be either y or \dot{y} and the Chebyshev polynomial coefficients of x are denoted by γ_k , we have the following conditions

$$x(-1) = \sum_{k=0}^{k=N} {}' \gamma_k T_k(-1) \quad (4.15)$$

$$x(1) = \sum_{k=0}^{k=N} {}' \gamma_k T_k(1) \quad (4.16)$$

leading to

$$\frac{1}{2}\gamma_0 - \gamma_1 + \gamma_2 + \cdots + (-1)^N \gamma_N = x(-1) \quad (4.17)$$

$$\frac{1}{2}\gamma_0 + \gamma_1 + \gamma_2 + \cdots + \gamma_N = x(1) \quad (4.18)$$

from which γ_0 and γ_1 can be recovered as

$$\gamma_0 = x(1) + x(-1) - 2(\gamma_2 + \gamma_4 + \gamma_6 + \cdots) \quad (4.19)$$

$$\gamma_1 = \frac{x(1) - x(-1)}{2} - (\gamma_3 + \gamma_5 + \gamma_7 + \cdots) \quad (4.20)$$

The zero coefficient for the other state is obtained by using the initial or final condition relationship between the state whose boundaries are constrained and the state whose boundaries are free. For example, if we know the boundary conditions for y , we have its Chebyshev expansion as

$$y(\tau) = \sum_{k=0}^{k=N} {}' \alpha_k T_k(\tau) \quad (4.21)$$

Taking the derivative leads to

$$\dot{y}(\tau) = \sum_{k=1}^{k=N} \alpha_k k \dot{T}_k(\tau) \quad (4.22)$$

Using the property that the derivative of the Chebyshev polynomial of the first kind is related with the Chebyshev polynomial of the second kind through

$$\frac{dT_k}{d\tau} = kU_{k-1} \quad (4.23)$$

we obtain

$$\dot{y}(\tau) = \sum_{k=1}^{k=N} \alpha_k k U_{k-1}(\tau) \quad (4.24)$$

where U_k denotes the Chebyshev polynomial of the second kind. Furthermore, using the following property of Chebyshev polynomials of the second kind

$$U_k(-1) = (k+1)(-1)^k \quad (4.25)$$

we obtain the initial condition constraint for \dot{y} as

$$\dot{y}(-1) = \sum_{k=1}^{k=N} \alpha_k k^2 (-1)^{(k+1)} \quad (4.26)$$

which is to be used to update the zero order coefficient of \dot{y} through Eq. 2.12.

The integration errors for this case are shown in Figs. 33 and 34. The boundary conditions for y are shown to have been satisfied to high accuracy.

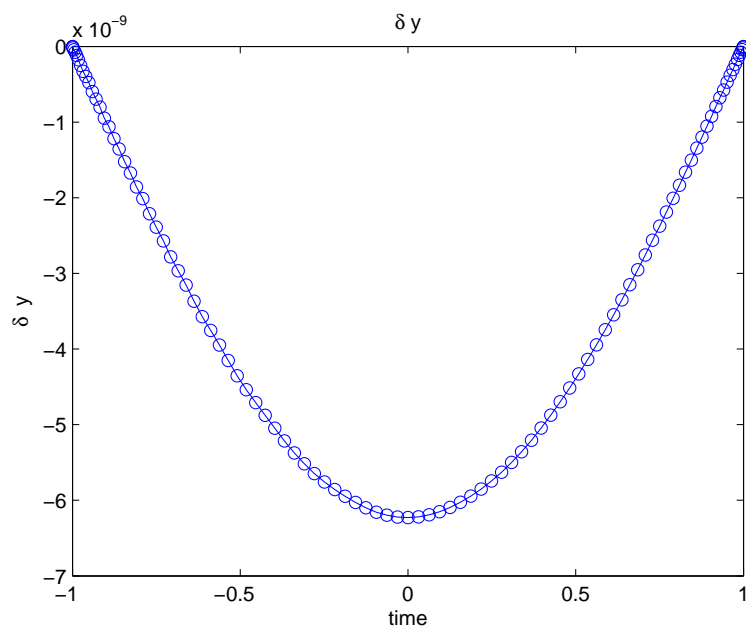


Fig. 33. Integration errors of y ($y(-1)$ and $y(1)$ are known)

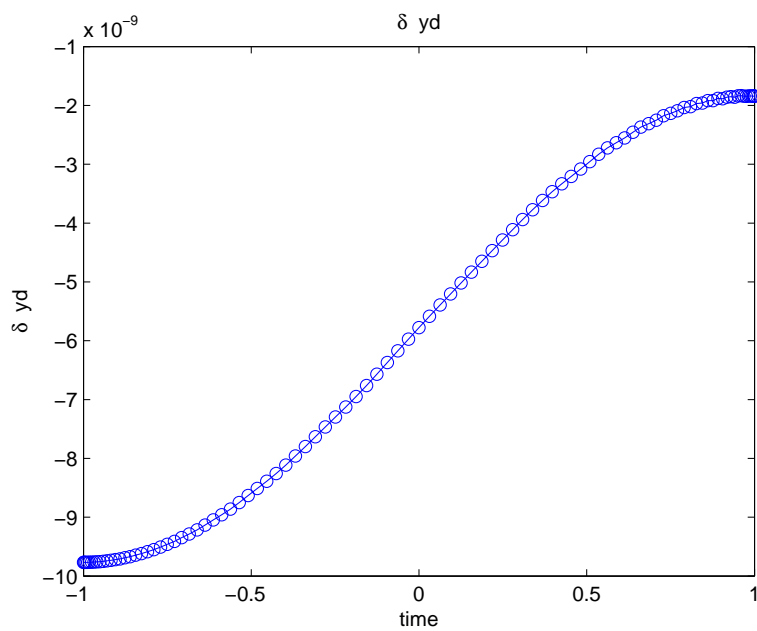


Fig. 34. Integration errors of \dot{y} ($y(-1)$ and $y(1)$ are known)

On the other hand, if we know $\dot{y}(-1)$, we can take the derivative of \dot{y} to get

$$\ddot{y}(-1) = -\lambda^2 y(-1) = \sum_{k=1}^{k=N} \beta_k k^2 (-1)^{k+1} \quad (4.27)$$

which leads to the initial condition for y as

$$y(-1) = - \sum_{k=1}^{k=N} \beta_k k^2 (-1)^{k+1} / \lambda^2 \quad (4.28)$$

The zero order coefficient of y can be obtained through

$$\alpha_0 = -2 \sum_{k=1}^{k=N} \beta_k k^2 (-1)^{k+1} / \lambda^2 - 2(-\alpha_1 + \alpha_2 + \cdots + (-1)^N \alpha_N) \quad (4.29)$$

The errors of the solutions for this case are shown in Figs. 35 and 36. Here the boundary conditions for \dot{y} are shown to have been satisfied to high accuracy.

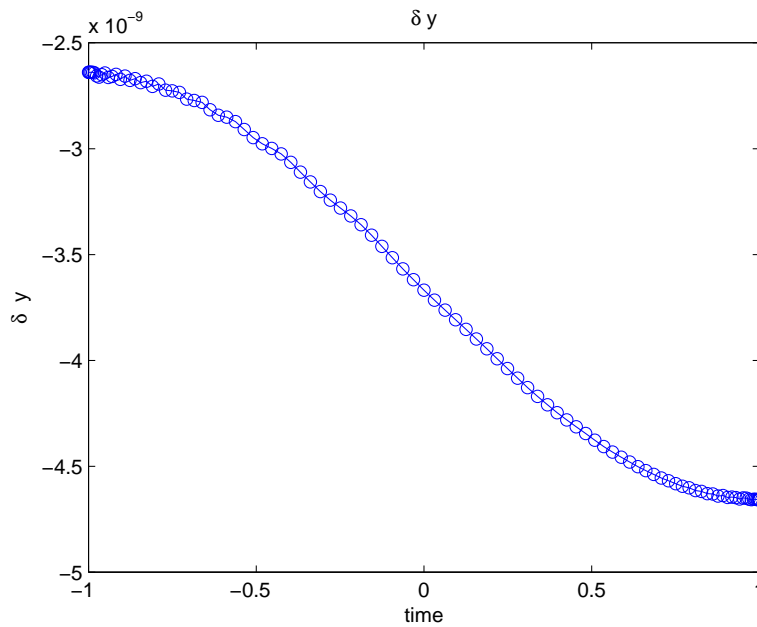


Fig. 35. Integration errors of y ($\dot{y}(-1)$ and $\dot{y}(1)$ are known)

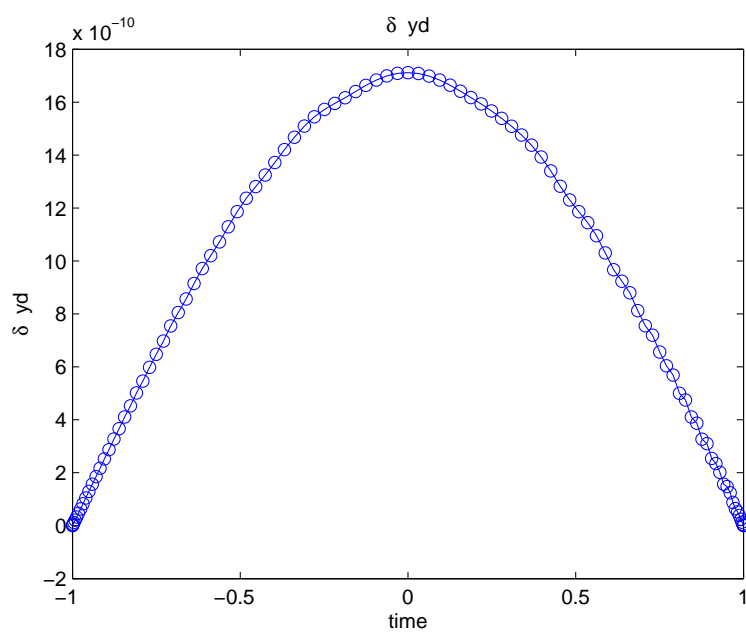


Fig. 36. Integration errors of \dot{y} ($\dot{y}(-1)$ and $\dot{y}(1)$ are known)

3. Solving BVPs of the Second and Third Kind

In this case, either $y(-1)$ and $\dot{y}(1)$ are known or $\dot{y}(-1)$ and $y(1)$ are known.

If we have $\dot{y}(1)$ and $y(-1)$, the zero coefficient conditions lead to the constraints

$$\alpha_0 = 2y(-1) - 2 \sum_{k=1}^{k=N} \alpha_k (-1)^k \quad (4.30)$$

$$\beta_0 = 2\dot{y}(1) - 2 \sum_{k=1}^{k=N} \beta_k \quad (4.31)$$

If we have $y(1)$ and $\dot{y}(-1)$, the zero coefficients lead to the constraints

$$\alpha_0 = 2y(1) - 2 \sum_{k=1}^{k=N} \alpha_k \quad (4.32)$$

$$\beta_0 = 2\dot{y}(-1) - 2 \sum_{k=1}^{k=N} \beta_k (-1)^k \quad (4.33)$$

Thus, both of these two types of problems can be solved essentially as a combination of an initial value problem for one state and a final value problem for another state. The two states are solved simultaneously using MCPI methods and the integration errors are shown in Figs. 37 through 40. Again, the boundary condition violations are negligible.

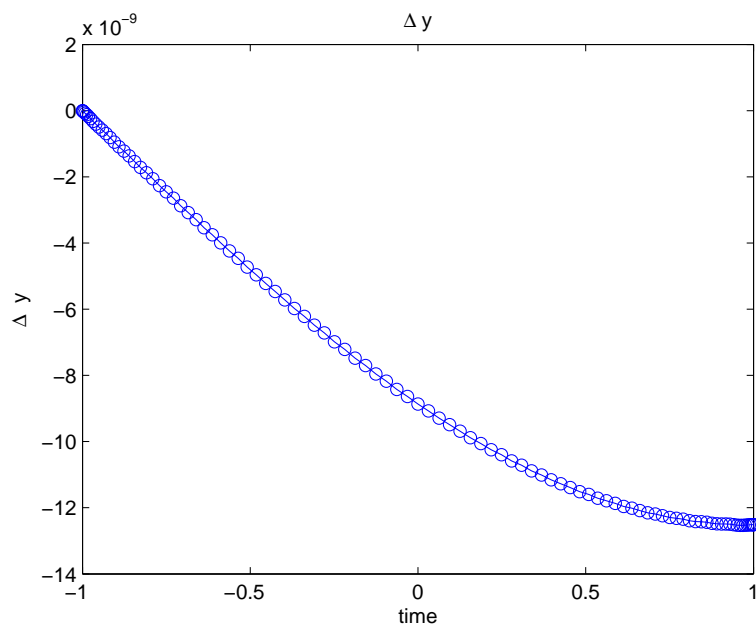


Fig. 37. Integration errors of y ($y(-1)$ and $yd(1)$ are known)

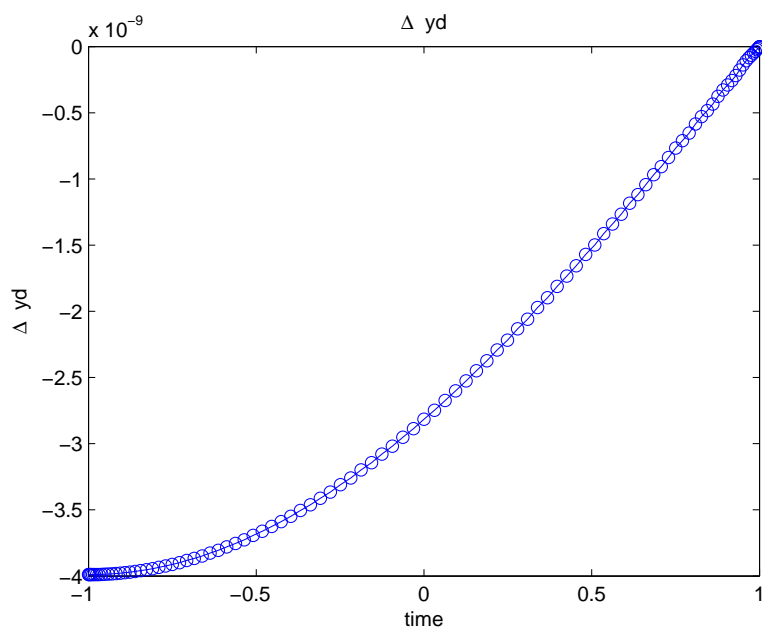


Fig. 38. Integration errors of \dot{y} ($y(-1)$ and $yd(1)$ are known)

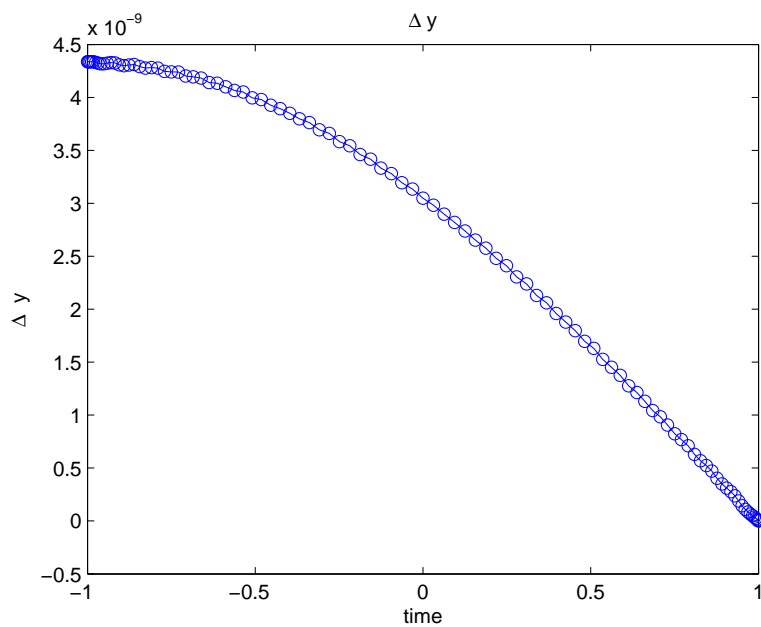


Fig. 39. Integration errors of y ($y(1)$ and $y(-1)$ are known)

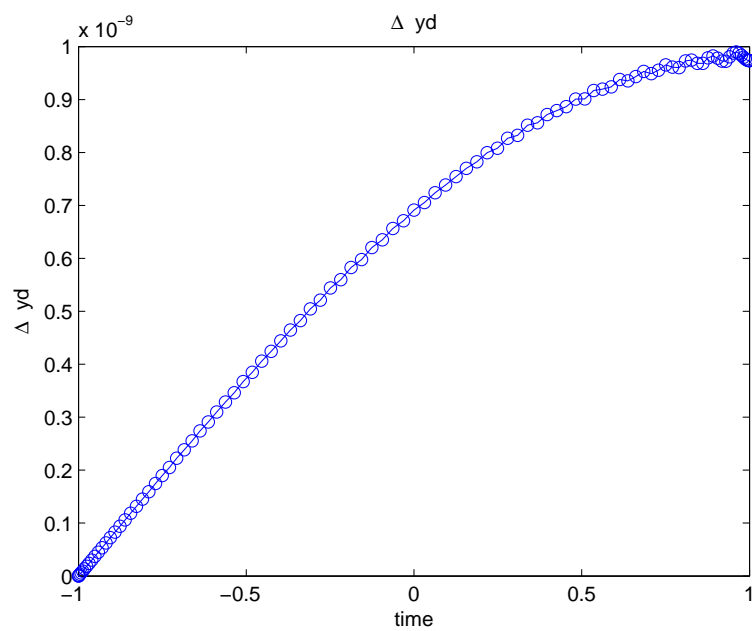


Fig. 40. Integration errors of \dot{y} ($y(1)$ and $y(-1)$ are known)

D. Applications to Lambert's Problem

The classical Lambert's problem is essentially a BVP of the first kind. Because of its significance in celestial mechanics fields, we present the details of using MCPI methods to solve this problem. The considered dynamic system satisfies Eqs. 3.39 to 3.41.

The time interval is $[0, T_f]$ and the boundary conditions of $x(0)$, $x(T_f)$, $y(0)$, $y(T_f)$, $z(0)$ and $z(T_f)$ are given. Reorganizing the second order ODE as a set of first order ODEs and normalizing the time to the range of $[-1, 1]$, one obtain

$$\frac{dx}{d\tau} = \frac{T_f}{2} \dot{x} \quad (4.34)$$

$$\frac{dy}{d\tau} = \frac{T_f}{2} \dot{y} \quad (4.35)$$

$$\frac{dz}{d\tau} = \frac{T_f}{2} \dot{z} \quad (4.36)$$

$$\frac{d\dot{x}}{d\tau} = -\frac{T_f}{2} \frac{\mu}{r^3} x \quad (4.37)$$

$$\frac{d\dot{y}}{d\tau} = -\frac{T_f}{2} \frac{\mu}{r^3} y \quad (4.38)$$

$$\frac{d\dot{z}}{d\tau} = -\frac{T_f}{2} \frac{\mu}{r^3} z \quad (4.39)$$

Assuming the six states are approximated by the Chebyshev polynomials

$$x(\tau) = \sum_{k=0}^{k=N} \alpha_k T_k(\tau) \quad (4.40)$$

$$y(\tau) = \sum_{k=0}^{k=N} \beta_k T_k(\tau) \quad (4.41)$$

$$z(\tau) = \sum_{k=0}^{k=N} \gamma_k T_k(\tau) \quad (4.42)$$

$$\dot{x}(\tau) = \sum_{k=0}^{k=N} \zeta_k T_k(\tau) \quad (4.43)$$

$$\dot{y}(\tau) = \sum_{k=0}^{k=N} \eta_k T_k(\tau) \quad (4.44)$$

$$\dot{z}(\tau) = \sum_{k=0}^{k=N} \xi_k T_k(\tau) \quad (4.45)$$

MCPI methods lead to

$$\vec{\alpha} = C_\alpha \frac{T_f}{2} \dot{\vec{x}} \quad (4.46)$$

$$\vec{\beta} = C_\alpha \frac{T_f}{2} \dot{\vec{y}} \quad (4.47)$$

$$\vec{\gamma} = C_\alpha \frac{T_f}{2} \dot{\vec{z}} \quad (4.48)$$

where $\vec{\alpha}$, $\vec{\beta}$, and $\vec{\gamma}$ are the coefficients written in the vector form; and $\dot{\vec{x}}$, $\dot{\vec{y}}$ and $\dot{\vec{z}}$ are the velocity evaluated at the CGL points.

The conditions on the initial and final positions require the zero and first order

coefficients to satisfy

$$\alpha_0 = x(0) + x(T_f) - 2(\alpha_2 + \alpha_4 + \alpha_6 + \dots) \quad (4.49)$$

$$\alpha_1 = \frac{x(0) - x(T_f)}{2} - (\alpha_3 + \alpha_5 + \alpha_7 + \dots) \quad (4.50)$$

$$\beta_0 = y(0) + y(T_f) - 2(\beta_2 + \beta_4 + \beta_6 + \dots) \quad (4.51)$$

$$\beta_1 = \frac{y(0) - y(T_f)}{2} - (\beta_3 + \beta_5 + \beta_7 + \dots) \quad (4.52)$$

$$\gamma_0 = z(0) + z(T_f) - 2(\gamma_2 + \gamma_4 + \gamma_6 + \dots) \quad (4.53)$$

$$\gamma_1 = \frac{z(0) - z(T_f)}{2} - (\gamma_3 + \gamma_5 + \gamma_7 + \dots) \quad (4.54)$$

The zero order coefficients for the velocities are computed as follows

$$\zeta_0 = \frac{2}{T_f} \sum_{k=1}^{k=N} \alpha_k k^2 (-1)^{(k+1)} \quad (4.55)$$

$$\eta_0 = \frac{2}{T_f} \sum_{k=1}^{k=N} \beta_k k^2 (-1)^{(k+1)} \quad (4.56)$$

$$\xi_0 = \frac{2}{T_f} \sum_{k=1}^{k=N} \gamma_k k^2 (-1)^{(k+1)} \quad (4.57)$$

In this way, the MCPI algorithm iterates to solve for the solutions of the positions and the velocities.

Now the satellite orbit propagation problem discussed in Section 2 of Chapter I is presented as a case study. The F and G solutions from that problem provide a baseline to verify the obtained results.

MCPI methods are compared with the elegant approach developed by Battin [41] as well as fsolve. Fsolve is a MATLAB nonlinear equation solver (for this case fsolve uses the Trust-Region Dogleg Method [48] as the nonlinear algorithm). It is important to keep in mind that Battin's approach can only solve the classical Lambert's problem if there is no perturbation involved except the inverse-squared gravity field, whereas both fsolve and MCPI methods have the capability to solve more general cases. The

initial conditions for fsolve are chosen to be close to the actual solutions. MCPI method starts the iteration by assuming zero velocity solutions and constant position solutions for all the three dimensions, where the constants are chosen as the initial positions.

The initial positions and final positions from the F and G solutions are the input variables to the three methods. The magnitude of the difference vector between the obtained initial velocity and the true initial velocity is defined as the error of the method. Figure 41 shows that MCPI method achieves a 50 to 80 speedup over the fsolve reference solution. Figures 42 and 43 show that when the MCPI method achieves close or better accuracy than the reference Battin's solution (Which has a specified convergence tolerance of 10^{-8}), the MCPI method is faster than the Battin's reference solution as long as the time interval is less than 15% of the orbit, although Battin's method is generally more efficient and applicable to arbitrary time of flight.

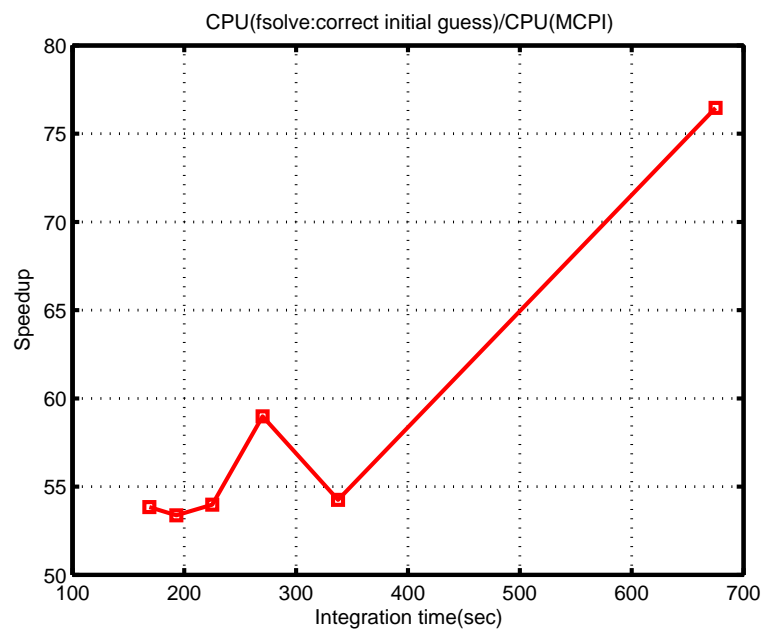


Fig. 41. Speedup of MCPI over the fsolve reference solution

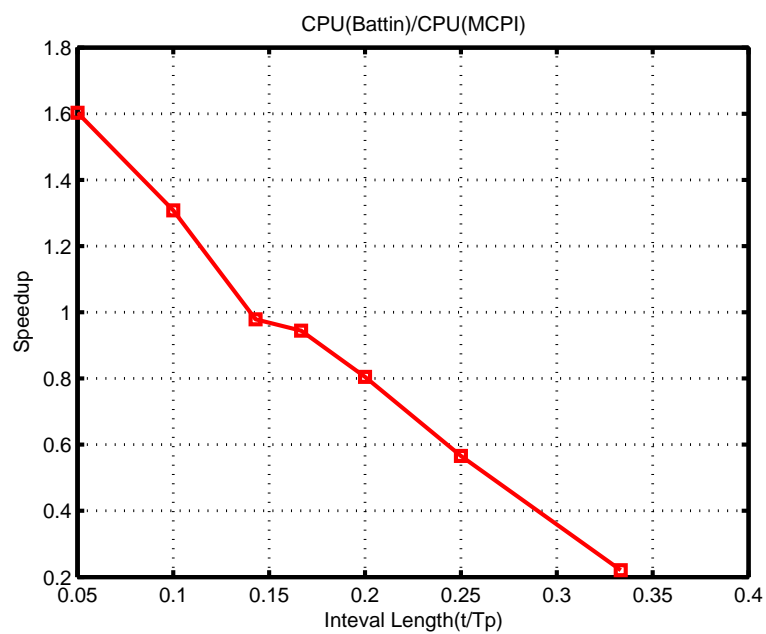


Fig. 42. Speedup of MCPI over a Battin's reference solution

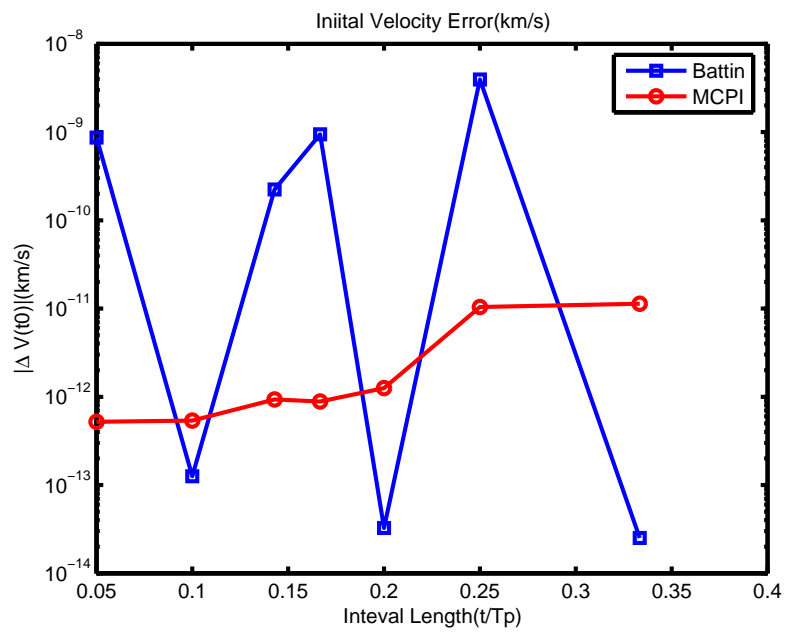


Fig. 43. Errors of MCPI and a Battin's reference solution

E. An Optimal Control Example: Earth to Apophis Optimal Trajectory Design

1. Introduction

The near-Earth asteroid 99942 Apophis has received a lot of attention since it was discovered on June 2004. Apophis will pass by the Earth on 2029 inside the geostationary orbit at about six times the Earth's radius from its geocenter. Following the 2029 encounter, another encounter is expected in 2036. The 2036 encounter remains highly uncertain and the current best estimate of Earth-Apophis impact probability on April 13, 2036 was calculated as 1 in 45000. Several researchers and groups have proposed to launch a small spacecraft to rendezvous with Apophis first [49, 50]. By using a beacon, transponder, or reflector on or near the asteroid, we can characterize its spin state, surface conditions, composition, etc. Using these information, we can further refine orbital propagation dynamic models for Apophis. Only with more accurate information on the orbit can wise decisions be made about whether a deflection mission is required or not.

A standard approach for an Earth to Apophis mission is to use a chemical rocket for launch to provide the initial needed Δv , which is denoted as V_1 in this section. Chemical thrusters/engines are also used for the rendezvous phase to provide the final rendezvous Δv , which is denoted as V_2 . The launch vehicle size and cost is dictated by the maximum V_1 they can provide for a prescribed mass leaving the vicinity of Earth. The rendezvous V_2 provided by the chemical thrusters also has some limitations. Large V_2 (for a given mass from Earth) means a large fraction of the mass must be propulsion system and fuel, decreasing the mass fraction of the scientific payload. In this traditional approach, designers usually generate “pork chop” plots that show V_1 and V_2 with respect to different launch date and time of flight [51]. The preferred launch dates/windows/and actual time of flight are picked mainly based on their

small V_1 and V_2 . Launch dates when the C_3 values (square of V_1) are beyond the launch vehicle capability or when the V_2 values are beyond the thrust capability must obviously be avoided. The sensitivities to a short fixed launch window cause problems in reality: either the spacecraft has to be built much earlier so that it can be launched in the ideal launch window, or a mission has to be delayed for a long time once it misses such a window.

Low thrust propulsion has been extensively studied for interplanetary applications with their successful applications. With the successful Deep Space 1 mission and the Dawn Spacecraft, it is now considered a feasible alternative to two-impulsive mission. Although the thrust magnitude is usually less than $0.1N$, low thrust propulsion systems usually have high specific impulse and high efficiency. However, the limitation of the thrust magnitude may not lead to a sufficiently short time of flight for the mission.

Junkins et al. have studied a hybrid impulsive with low thrust propulsion strategy for an Earth to Apophis mission [52, 53]. For this approach, low thrust propulsion is used during the flight and impulses are available at both the launch phase and the rendezvous phase. For a prescribed thrust magnitude, the thrust steering angle is sought that minimize the norm of the terminal velocity impulses. The authors use a Newton method that incorporates the state parameter transition tensors to solve the two-point boundary value problems. Using the implicit function theorem, they generate an extremal field map family of optimal solutions with variations of launch date, time of flight and thrust magnitude. Bai et al. applied a novel homotopy method to an Earth to Apophis mission analysis and further studied the hybrid impulsive and low thrust propulsion strategy [54]. They showed that using a $0.05N$ constant thrust and choosing the time of flight from 250 days to 300 days allow the launch window for the Earth to Apophis mission to span a full year from April 2012 to April 2013 for a

standard launch vehicle. Due to the significance of this discovery, the same problem is studied here again using the proposed MCPI methods.

2. Frame Definition and Dynamic Equations

A two-dimensional Earth to Apophis trajectory is considered. This simplification is valid for a preliminary study because the inclination of Apophis is about three degrees. The spacecraft is assumed to be a point mass with a variable mass m . As shown in Fig. 44, the position of the spacecraft in the solar-centric polar coordinates is (r, θ) , where r is the distance of the spacecraft to the Sun and θ is the phase angle with respect to some inertial axis. More, u is the velocity along the radial direction and v is the velocity along the local horizontal direction. The angle between the thrust direction and the local horizontal direction is the control variable represented by β .

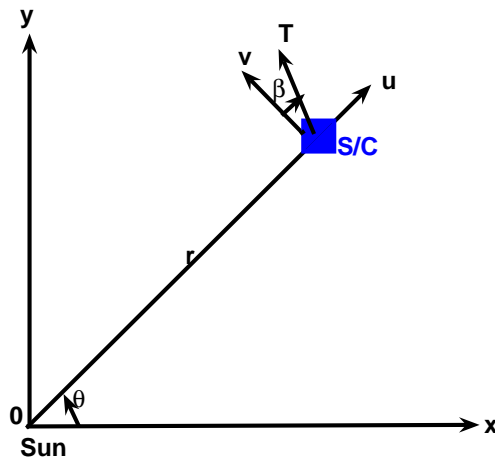


Fig. 44. Frame, State and Control Definition

The dynamic equations for the spacecraft are

$$\dot{r} = u \quad (4.58)$$

$$\dot{\theta} = v/r \quad (4.59)$$

$$\dot{u} = v^2/r - \mu/r^2 + T/m \sin \beta \quad (4.60)$$

$$\dot{v} = -uv/r + T/m \cos \beta \quad (4.61)$$

where T is the constant thrust.

The mass flow rate α is constant and the resulting mass equation is

$$m = m_0 - \alpha(t - t_0), t_0 \leq t \leq t_f \quad (4.62)$$

where m_0 is the initial mass, t_0 is the initial time, and t_f is the final time. The thrust magnitude is related to the mass flow rate through

$$T = c\alpha \quad (4.63)$$

where we choose $c = 1.872$ as suggested by Oberle and Taubert. [55]

All the values of length are nondimensionalized by 1.496×10^{11} meters, which is the Earth's distance to the Sun. All the values of time are nondimensionalized by 5.023×10^6 sec. Subsequently, μ in Eq. 4.60 becomes one after the nondimensionalization. The initial mass is nondimensionalized as $m_0 = 1$. We choose the mass flow rate to be $\alpha = 0.007487$. The resulting nondimensionalized thrust is 0.014, which corresponds to 0.05N for a 600kg spacecraft. The position and velocity of the Earth on the launch date and the position and velocity of Apophis on the rendezvous date are the ephemeris data obtained through Jet Propulsion Laboratory online Horizon system.

3. Optimality Criterion and First Order Conditions

Motivated by Junkins et al. [52, 53], the goal of the optimal control is to minimize the total impulse energy at the terminal times, defined as

$$J = \frac{1}{2}V_1^2 + \frac{1}{2}V_2^2 \quad (4.64)$$

The control variable is the thrust steering angle β and the optimization constraints are the dynamic equations from Eq. 4.58 to Eq. 4.62.

For the polar reference frame defined in Figure 44, we have

$$V_1^2 = (u_{s/c}(t_0) - u_{Apophis}(t_0))^2 + (v_{s/c}(t_0) - v_{Apophis}(t_0))^2 \quad (4.65)$$

$$V_2^2 = (u_{s/c}(t_f) - u_{Apophis}(t_f))^2 + (v_{s/c}(t_f) - v_{Apophis}(t_f))^2 \quad (4.66)$$

The Hamiltonian is defined as

$$H = \lambda_r u + \lambda_\theta r/v + \lambda_u (v^2/r - \mu/r^2 + T/m \sin \beta) + \lambda_v (-uv/r + T/m \cos \beta) \quad (4.67)$$

The co-state equations obtained through Pontryagin's principle are

$$\dot{\lambda}_r = -\lambda_u (-v^2/r^2 + 2/r^3) - \lambda_v uv/r^2 + \lambda_\theta v/r^2 \quad (4.68)$$

$$\dot{\lambda}_\theta = 0 \quad (4.69)$$

$$\dot{\lambda}_u = -\lambda_r + \lambda_v v/r \quad (4.70)$$

$$\dot{\lambda}_v = -2\lambda_u v/r + \lambda_v u/r - \lambda_\theta/r; \quad (4.71)$$

and the optimal thrust direction is

$$\beta = \text{atan2}(-\lambda_u, -\lambda_v) \quad (4.72)$$

where $\text{atan2}()$ is the four-quadrant inverse tangent function.

The transversality conditions lead to the following two-point boundary conditions

$$r(t_0) = r_{Earth}(t_0) \tag{4.73}$$

$$\theta(t_0) = \theta_{Earth}(t_0) \tag{4.74}$$

$$r(t_f) = r_{Apophis}(t_f) \tag{4.75}$$

$$\theta(t_f) = \theta_{Apophis}(t_f) \tag{4.76}$$

$$\lambda_u(t_0) = -(u(t_0) - u_{Earth}(t_0)) \tag{4.77}$$

$$\lambda_v(t_0) = -(v(t_0) - v_{Earth}(t_0)) \tag{4.78}$$

$$\lambda_u(t_f) = u(t_f) - u_{Apophis}(t_f) \tag{4.79}$$

$$\lambda_v(t_f) = v(t_f) - v_{Apophis}(t_f) \tag{4.80}$$

4. Solutions from MCPI Methods

The boundary conditions for the six states and six co-states are constrained by the following equations during the MCPI iterations

$$\alpha_0^r = r(t_0) + r(t_f) - 2(\alpha_2^r + \alpha_4^r + \alpha_6^r + \dots) \quad (4.81)$$

$$\alpha_1^r = \frac{r(t_f) - r(t_0)}{2} - (\alpha_3^r + \alpha_5^r + \alpha_7^r + \dots) \quad (4.82)$$

$$\alpha_0^\theta = \theta(t_0) + \theta(t_f) - 2(\alpha_2^\theta + \alpha_4^\theta + \alpha_6^\theta + \dots) \quad (4.83)$$

$$\alpha_1^\theta = \frac{\theta(t_f) - \theta(t_0)}{2} - (\alpha_3^\theta + \alpha_5^\theta + \alpha_7^\theta + \dots) \quad (4.84)$$

$$u(t_0) = \frac{2}{t_f} \sum_{k=1}^{k=N} \alpha_k^r k^2 (-1)^{(k+1)} \quad (4.85)$$

$$v(t_0) = \frac{2}{t_f} r(t_0) \sum_{k=1}^{k=N} \alpha_k^\theta k^2 \quad (4.86)$$

$$\alpha_0^{\lambda_u} = \lambda_u(t_0) + \lambda_u(t_f) - 2(\alpha_2^{\lambda_u} + \alpha_4^{\lambda_u} + \alpha_6^{\lambda_u} + \dots) \quad (4.87)$$

$$\alpha_1^{\lambda_u} = \frac{\lambda_u(t_f) - \lambda_u(t_0)}{2} - (\alpha_3^{\lambda_u} + \alpha_5^{\lambda_u} + \alpha_7^{\lambda_u} + \dots) \quad (4.88)$$

$$\alpha_0^{\lambda_v} = \lambda_v(t_0) + \lambda_v(t_f) - 2(\alpha_2^{\lambda_v} + \alpha_4^{\lambda_v} + \alpha_6^{\lambda_v} + \dots) \quad (4.89)$$

$$\alpha_1^{\lambda_v} = \frac{\lambda_v(t_f) - \lambda_v(t_0)}{2} - (\alpha_3^{\lambda_v} + \alpha_5^{\lambda_v} + \alpha_7^{\lambda_v} + \dots) \quad (4.90)$$

$$\lambda_r(t_0) = -\frac{2}{t_f} r(t_0) \sum_{k=1}^{k=N} \alpha_k^{\lambda_u} k^2 + \lambda_v(t_0) v(t_0) / r(t_0) \quad (4.91)$$

$$\lambda_\theta(t_0) = r(t_0) (-2\lambda_u(t_0) v(t_0) / r(t_0) + \lambda_v(t_0) u(t_0) / (r(t_0) - \frac{2}{t_f} \sum_{k=1}^{k=N} \alpha_k^{\lambda_v} k^2 (-1)^{(k+1)})) \quad (4.92)$$

In the simulation, we assume that the spacecraft leaves Earth on April 18, 2012 and rendezvous with Apophis on August 6, 2012. The time of flight is 110 days.

The homotopy method presented in Ref. [54] is first use to find a solution. Because the boundary condition violations in Eqs. 4.73 to 4.80 through the homotopy approach are of the magnitude of 10^{-14} , we use this highly accurate solution to ver-

ify the results from the MCPI method. Thus all the errors shown in the figures are based on the solutions from the homotopy approach. The homotopy method takes about six seconds to find the solution when a very close initial guess is used. For reference, we have also used SNOPT [56] as a nonlinear programming solver to solve the same problem and the computation time is about three seconds in MATLAB. The boundary condition violations using SNOPT are in the magnitude of 10^{-11} . The MCPI method uses 192 iterations and finds the solution in only 0.15 seconds with the boundary condition violations in the magnitude of 10^{-14} . For the MCPI method to start, the solutions of the radius and phase angle are chosen as straight lines by using the initial conditions and the final conditions; the solution of the radial velocity is chosen as a zero vector; the solution of the tangential velocity is obtained from the radius information by assuming a circular orbit; and all the co-states are chosen as zero vectors. However, various test cases show that the MCPI method is not sensitive to these initial guesses.

The optimal transfer orbit is shown in Fig. 45 and the resulting optimal thrust angles are shown in Fig. 46. Figures 47 through 51 show the relative errors of the radius, phase angle, radial velocity, tangential velocity, and thrust angle. Again these results show that MCPI method obtains highly accurate solutions and the boundary condition violations are always negligible. We mention, however, when we increase the time of flight, that the MCPI method experiences convergence difficulties. Therefore, for BVPs, the MCPI method's convergence for long time intervals is the biggest drawback. In contrast to the case for MCPI solutions of IVPs, we have not derived a general MCPI algorithm with sub-division of the time of flight interval.

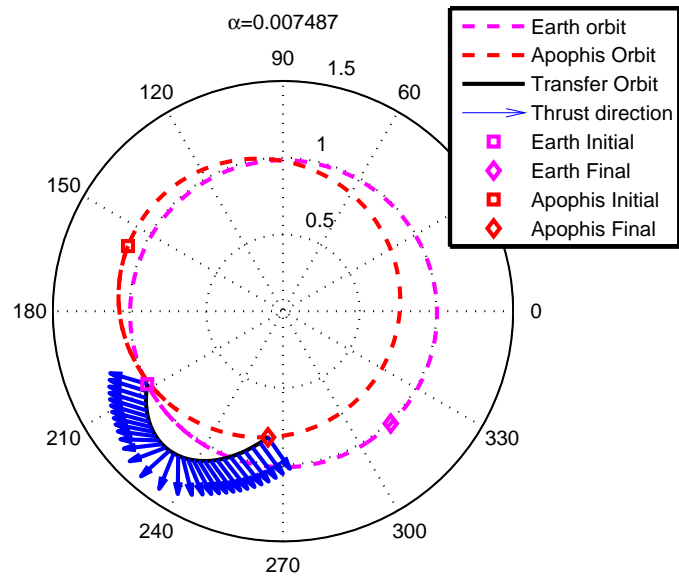


Fig. 45. Transfer orbit (Earth to Apophis)

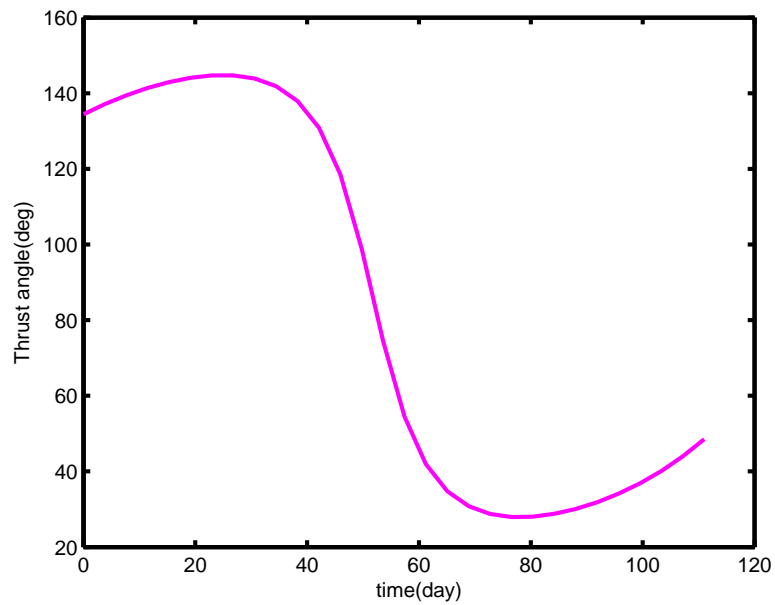


Fig. 46. Thrust angle history (Earth to Apophis)

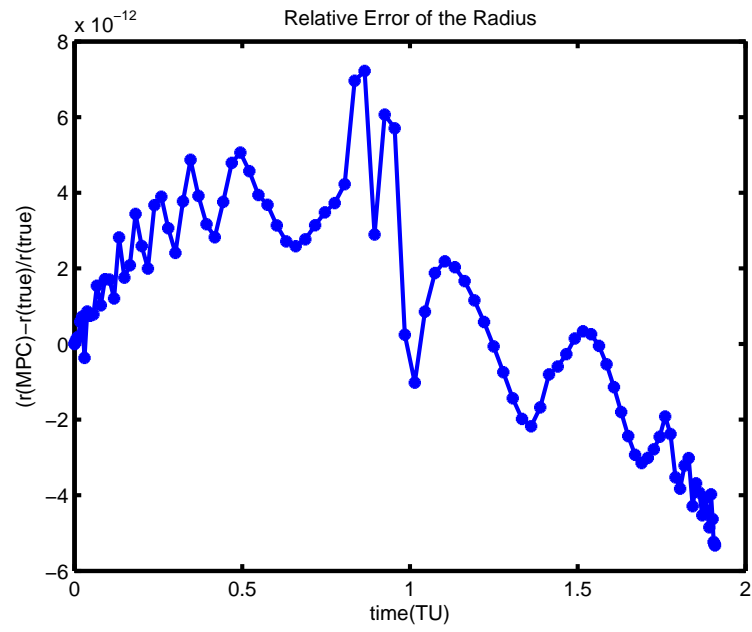


Fig. 47. Relative errors of the radius (Earth to Apophis)

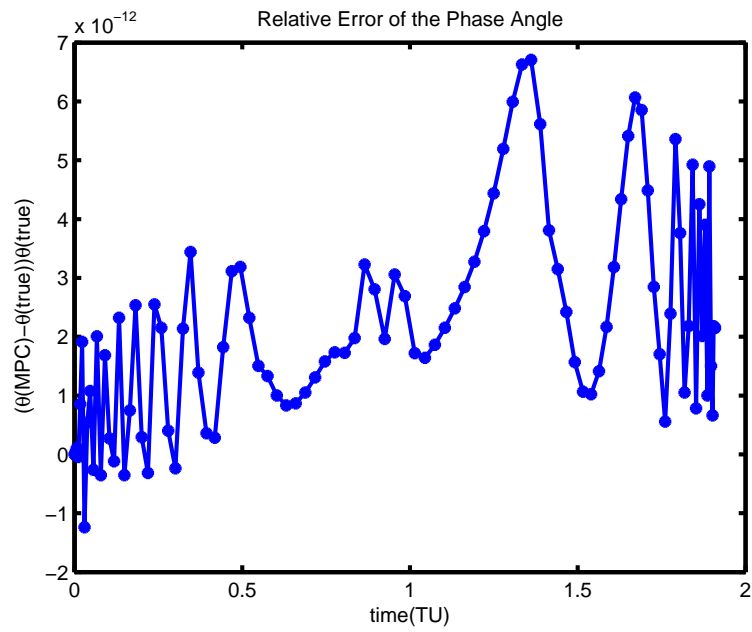


Fig. 48. Relative errors of the phase angle (Earth to Apophis)

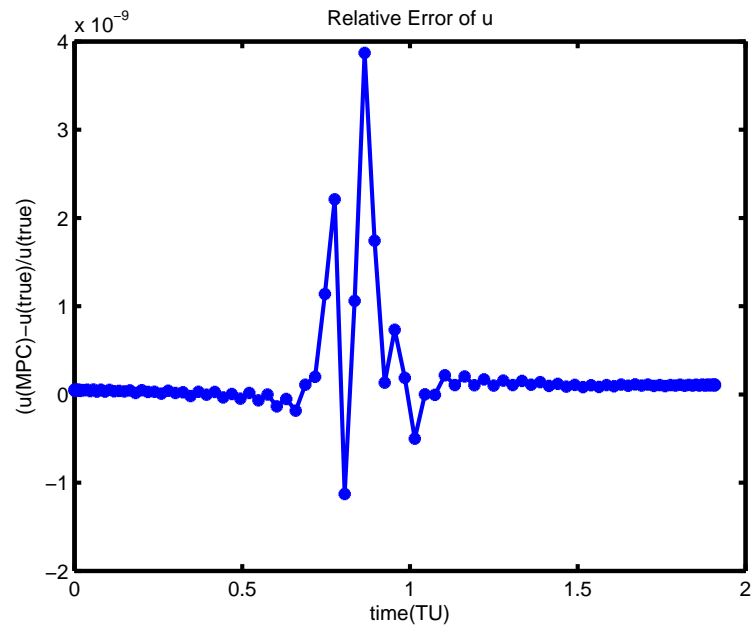


Fig. 49. Relative errors of the radial velocity (Earth to Apophis)

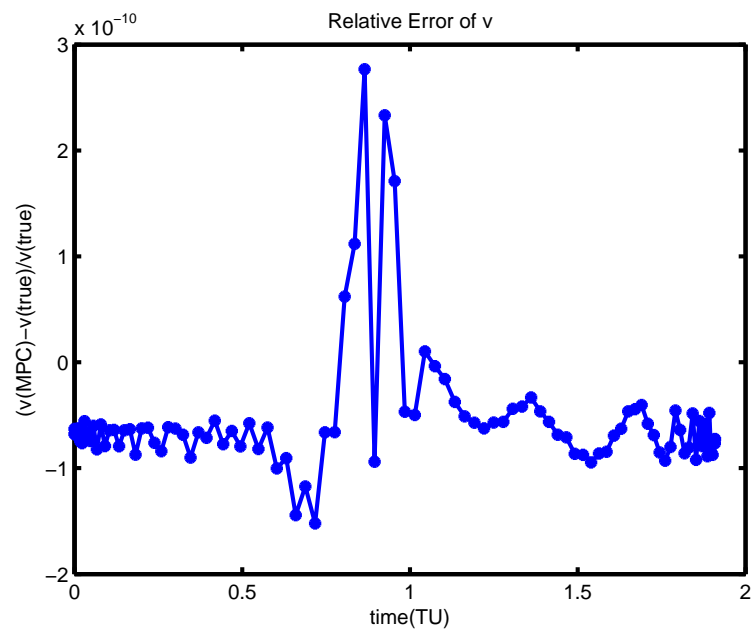


Fig. 50. Relative errors of the tangential velocity (Earth to Apophis)

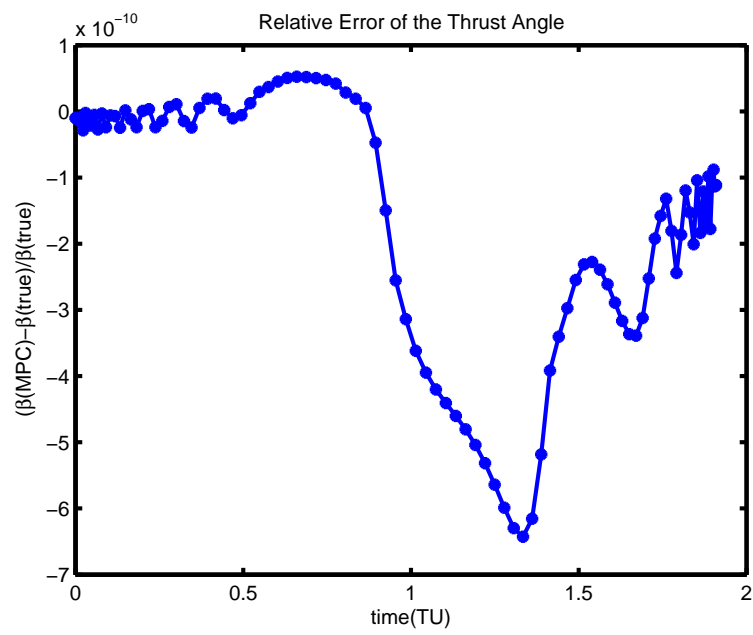


Fig. 51. Relative errors of the thrust angle (Earth to Apophis)

5. Solutions from a Chebyshev Pseudospectral Method

Because of the recent successful applications of the pseudospectral methods [29, 57, 58], we solve the same Earth to Apophis trajectory problem here using a pseudospectral method in order to compare its performance with that of MCPI methods. To be consistent, a Chebyshev pseudospectral method [29] is chosen, such that both the MCPI method and the Chebyshev pseudospectral method use the Chebyshev polynomials of order 100 and use the same nodes to approximate the solutions. The difference between these two approaches is fundamental. The pseudospectral method transforms the original problem to a nonlinear programming problem, thus it belongs to the category of the direct approach. However, the proposed MCPI method transforms the original problem to a two-point BVP by using Pontryagin’s principle, thus it is a variation of the indirect approach. However, in contrast to most existing indirect methods, gradient information is not required for MCPI methods.

It turns out that the initial guess for the Chebyshev pseudospectral method requires user insight or some preliminary trial and error process. To circumvent this issue, we provide an initial condition that is beneficial for the Chebyshev pseudospectral method. The starting guess for the radius, phase angle, radial velocity, and tangential velocity is chosen as the same starting guess used for the MCPI method. However, the initial guess of the steering angle that is required for the Chebyshev pseudospectral method is chosen as the true solution from the homotopy method. We note that this information would not generally be available for the pseudospectral method, so some other starting guess would be required. The Chebyshev pseudospectral method takes about 4.5 seconds to converge, where “fmincon” in MATLAB is used as the nonlinear optimizer. Thus the speedup of the MCPI method over the Chebyshev pseudospectral method is about 30 for this application. The termination

tolerance on the function value, tolerance on the constraint violation, and termination tolerance on the states for “fmincon” are set as 10^{-10} (“fmincon” does not converge when setting these parameters lower than these). Before looking at the results, we mention that all high dimensional nonlinear programming problems suffer, to a varying degree, from the curse of dimensionality. Seeking a minimum in a space of dimension 100 poses a computational challenge to achieve a high precision convergence to the actual minimum.

Figures 52 through 56 show the relative errors of the radius, phase angle, radial velocity, tangential velocity, and thrust angle. Comparing with Figs. 47 through 49, we can see that the solution of the states from the MCPI method has six to seven orders of magnitude better accuracy than the Chebyshev pseudospectral method. Figures 51 and 56 show that the optimal steering angle obtained from the MCPI method has about eight digits better accuracy than the Chebyshev pseudospectral method. The low accuracy of the Chebyshev pseudospectral method is also a consequence of the convergence characteristics of the nonlinear programming solver “fmincon”. We also experimented to use the correct histories of the states and steering angle from the homotopy method as the initial guess for the Chebyshev pseudospectral method. “fmincon” takes about one second to find the solutions and results the relative errors of the states less than 10^{-9} and the error of the control less than 10^{-7} . Consistent with the above comments, the nonlinear programming solver “fmincon” has difficulty isolating the final solution, although it converged efficiently to the approximate neighborhood of the solution. Utilizing Pontryagin’s principle and introducing the costate equations are the fundamental reasons for the MCPI method to achieve high accuracy. The computationally efficient structure of the MCPI method is the reason for both the high accuracy and significant speedup. However, as we pointed out earlier, the convergence of the MCPI method for long time intervals is not guaranteed. Thus the

Chebyshev pseudospectral method can solve more general optimal control problems than the MCPI method, at least at the current stage of our research.

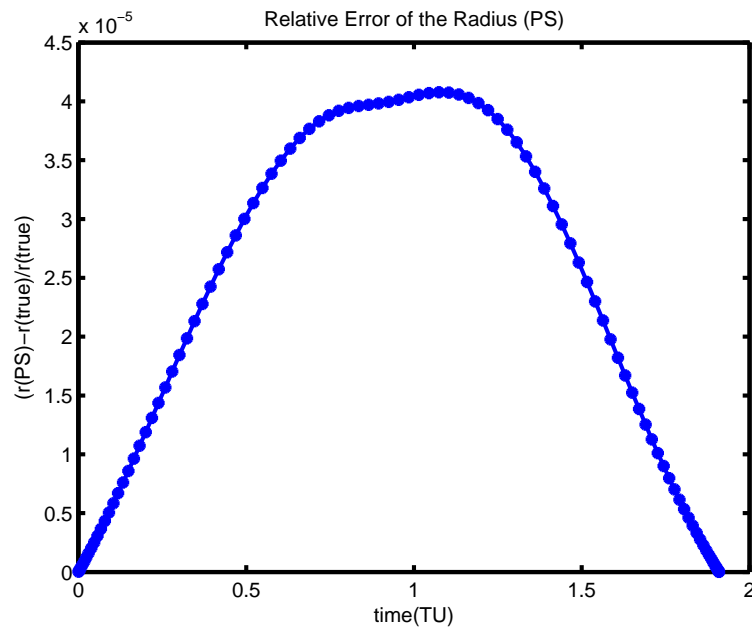


Fig. 52. Relative errors of the radius (pseudospectral)

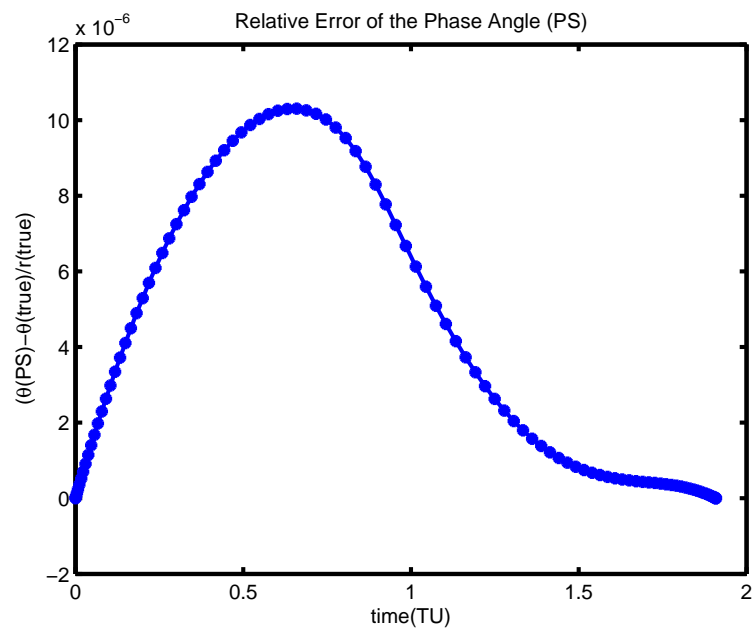


Fig. 53. Relative errors of the phase angle (pseudospectral)

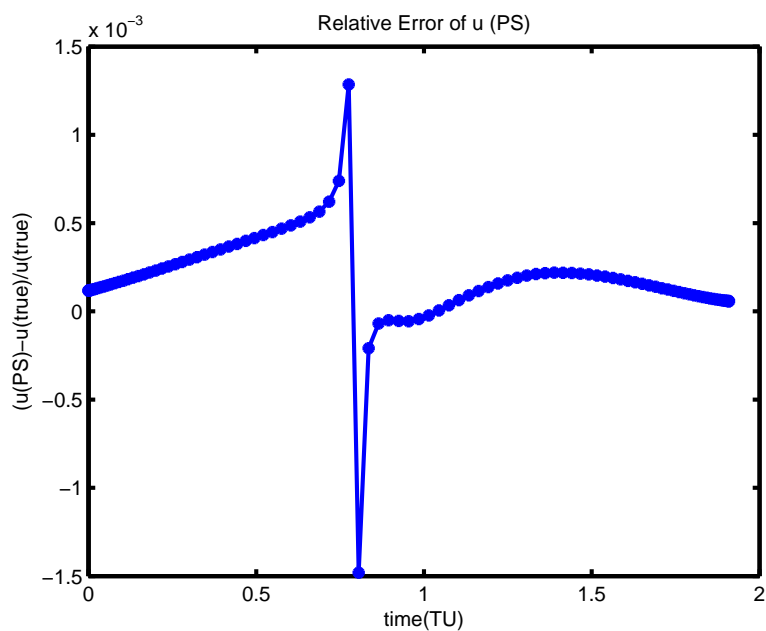


Fig. 54. Relative errors of the radial velocity (pseudospectral)

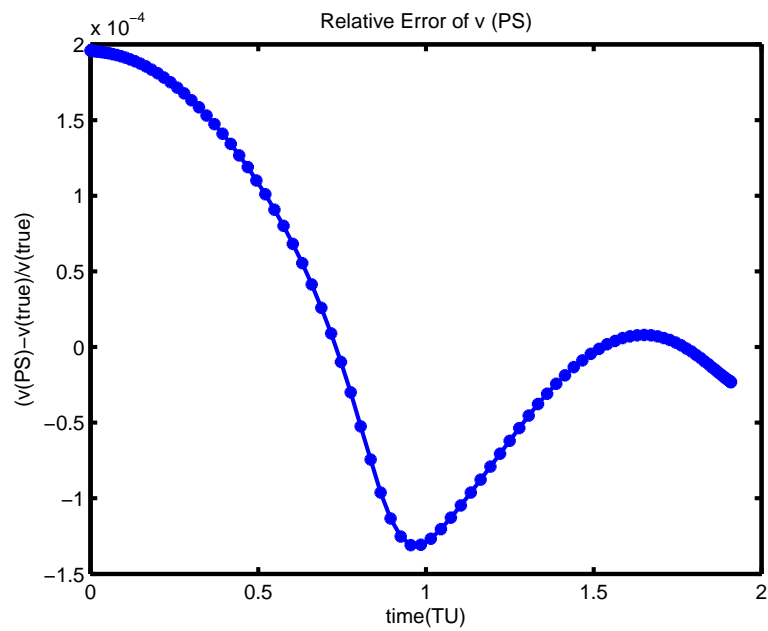


Fig. 55. Relative errors of the tangential velocity (pseudospectral)

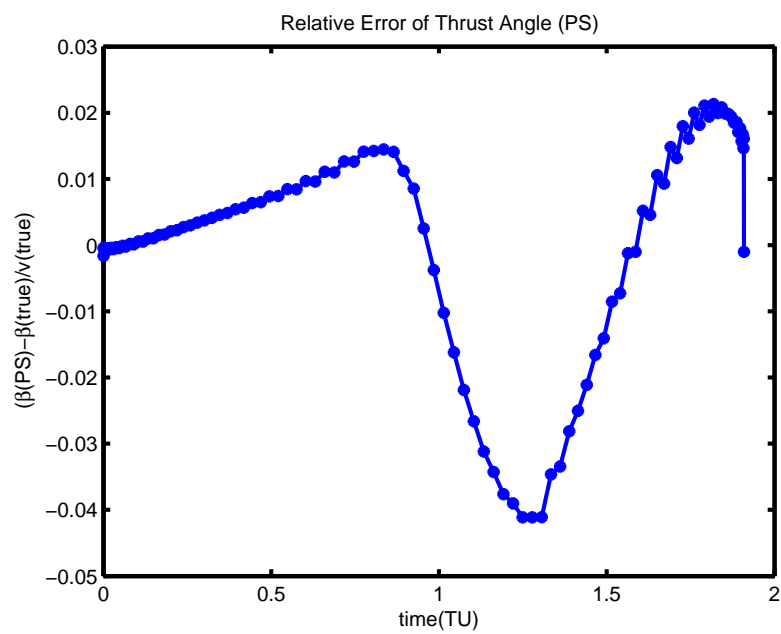


Fig. 56. Relative errors of the thrust angle (pseudospectral)

F. Summary

Through a detailed discussion on the process of using MCPI methods to solve a second order problem, this chapter illustrates that the proposed MCPI methods are applicable to solve general BVPs while not limited to BVPs of the first kind. For Lambert's problem, MCPI methods achieve both large speedup and high accuracy when compared with other methods, with the additional capability to solve perturbation-involved general cases. MCPI methods can also solve optimal control problems once the original problem has been transformed to a two-point boundary value problem. Both states and co-states are solved simultaneously. The results from the numerical examples presented in this chapter illustrate two significant advantages of using MCPI methods to solve BVPs. One is that the boundary conditions are always satisfied to high accuracy. Another is that the method is highly computationally efficient. However, the convergence of MCPI methods for solving BVPs, especially with long intervals, is not guaranteed. We discuss several techniques in the next chapter that will improve the convergence domain for MCPI methods.

CHAPTER V

DIFFERENT TECHNIQUES TO ENLARGE THE CONVERGENCE DOMAIN
OF MCPI METHODS

A. Introduction

In Chapter III, we show that the convergence domain for MCPI methods is restricted. Although the piecewise approach introduced in Chapter III extends MCPI methods to solve IVPs on an arbitrary time domain, the convergence of MCPI methods is not guaranteed for BVPs on a large domain, and a general approach to introduce a piecewise MCPI methods for BVPs have not been developed. Among the existing nonlinear transformation techniques that may either help the divergent iterations to converge or help the convergent iterations to converge faster than the original iterations [59, 60], Aitken's process is introduced as the first technique in this chapter to improve the convergence of MCPI methods. We present a linearization approach as a second technique. Inspired by the concept of proportional control, this chapter presents a correction control MCPI methodology as a third technique. The three techniques are applied for solving IVPs and BVPs in the numerical example section and their respective power is illustrated.

B. Aitken's Process

Aitken first proposed a process to accelerate the convergence rate of the sequences to find the greatest root of an algebraic equation using Bernoulli's method [61]. Consider

a sequence $\{x_n\}$ that converges to x in a geometric fashion as

$$x_n - x \approx K(x_{n-1} - x) \quad (5.1)$$

$$x_{n+1} - x \approx K(x_n - x) \quad (5.2)$$

where K is some constant called convergent rate. x can be solved using Eqs. 5.1 and 5.2 by eliminating the unknown constant K , leading to the form

$$x \approx \frac{x_{n-1}x_{n+1} - x_n^2}{x_{n+1} - 2x_n + x_{n-1}} \quad (5.3)$$

Aitken's process generates a new sequence $\{A_n\}$ from the old sequence $\{x_n\}$ through

$$A_n = \frac{x_{n-1}x_{n+1} - x_n^2}{x_{n+1} - 2x_n + x_{n-1}} \quad (5.4)$$

The second derived sequence is generated in a similar way

$$B_n = \frac{A_{n-1}A_{n+1} - A_n^2}{A_{n+1} - 2A_n + A_{n-1}} \quad (5.5)$$

and has been proved to converge faster in many applications than the sequence $\{A_n\}$ [59, 60]. Higher order sequences can be generated successively. However, as pointed by Aitken [61], it is more effective to use the derived sequences at an early stage rather than calculate all the lower order sequences and then formulate the higher order sequences.

Equation 5.4 can be written in another form as

$$A_n = x_{n+1} - \frac{(x_{n+1} - x_n)^2}{x_{n+1} - 2x_n + x_{n-1}} \quad (5.6)$$

This equation is numerically more stable than Eq. 5.4. Since as x_{n-1} , x_n , and x_{n+1} are close to x , both the denominator and numerator in Eq. 5.4 are close to zero, which will cause a first order cancelation; in such cases, the second part of Eq. 5.6

will also have both the denominator and numerator close to zero, but this is a second order cancelation that is numerically more stable [62]. Notice that Aitken's process is a special case of Shanks transformation [63]. For the sequence $\{x_n\}$, the k^{th} order Shanks transformation is defined by

$$S_{k,n} = \frac{\begin{vmatrix} x_{n-k} & \cdots & x_{n-1} & x_n \\ \Delta x_{n-k} & \cdots & \Delta x_{n-1} & \Delta x_n \\ \vdots & \vdots & \vdots & \vdots \\ \Delta x_{n-1} & \cdots & \cdots & \Delta x_{n+k-1} \end{vmatrix}}{\begin{vmatrix} 1 & 1 & 1 & 1 \\ \Delta x_{n-k} & \cdots & \Delta x_{n-1} & \Delta x_n \\ \vdots & \vdots & \vdots & \vdots \\ \Delta x_{n-1} & \cdots & \cdots & \Delta x_{n+k-1} \end{vmatrix}} \quad (5.7)$$

where $\Delta x_n = x_{n+1} - x_n$. Although in most cases, Shanks transformation is used for the divergent or slowly convergent sequence that has a partial summation form, Eq. 5.7 is a generalization of Aitken's process. For the first order case when $k = 1$, Shanks transformation reduces to Aitken's process. In this dissertation, to distinguish it from Shanks transformation, Aitken's process means the derived Aitken's process that is shown in Fig. 57.

We use a root solving problem to illustrate the power of these transformations for speeding up convergence. Consider the algebraic equation

$$x = e^{-x} \quad (5.8)$$

and assume the root can be solved iteratively through

$$x_n = e^{-x_{n-1}} \quad (5.9)$$

Figure 58 shows the convergence history of the original series as well as the series utilizing the first and second order Shanks transformation and the derived Aitken's process. These results confirm that although the first and second order series are derived from the original series, they converge much faster than the original series. The advantage of using Aitken's process is also shown to be significant. For Shanks transformation, higher orders lead to faster convergence.

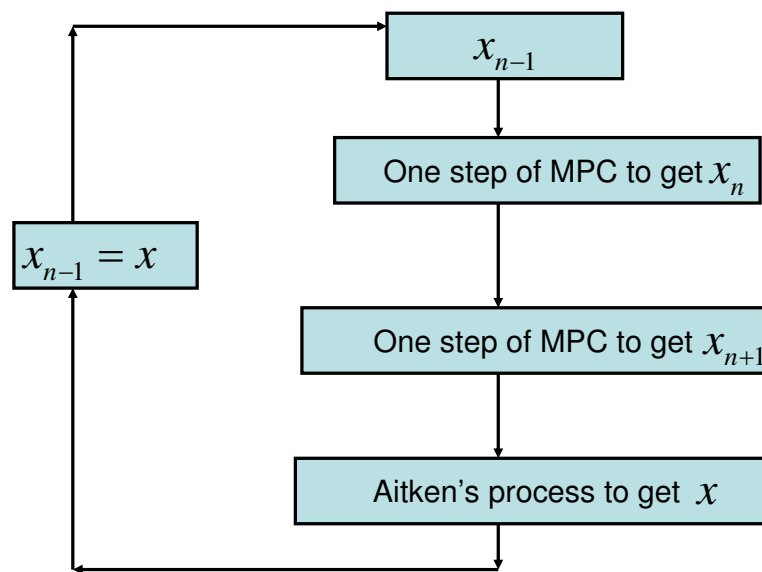


Fig. 57. Implement the derived Aitken's process

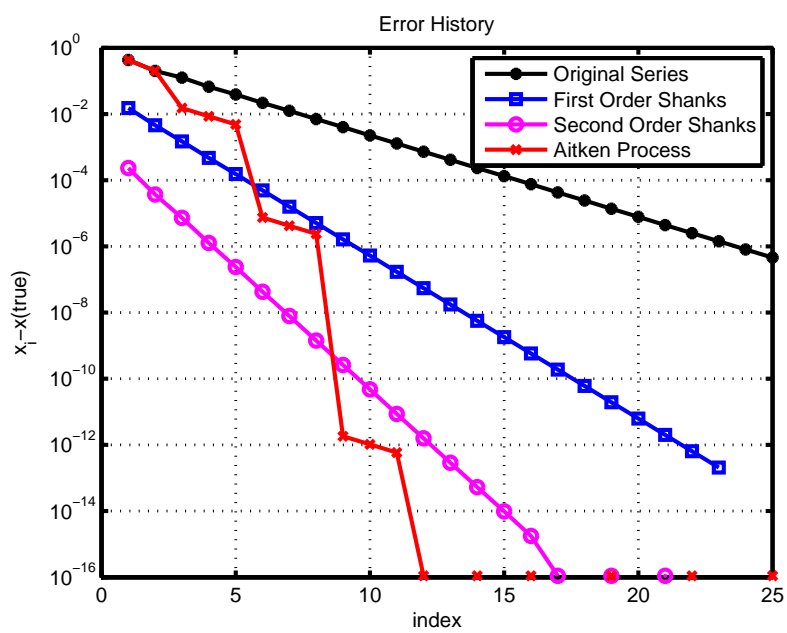


Fig. 58. Convergence history using different transformations

C. The Linearization Approach

For a first order differential equation

$$\frac{d\mathbf{x}}{d\tau} = \mathbf{f}(\tau, \mathbf{x}), -1 \leq \tau \leq 1 \quad (5.10)$$

the linearization approach generates a sequence of solutions by using the first-order Taylor expansion of function \mathbf{f} . The solution at the i^{th} iteration $\mathbf{x}^i(\tau)$ is related to the solution at the $(i-1)^{\text{th}}$ iteration $\mathbf{x}^{i-1}(\tau)$ through the formula

$$\mathbf{x}^i(\tau) = \mathbf{x}(-1) + \int_{-1}^{\tau} \left(\mathbf{f}(\mathbf{x}^{i-1}, s) + \frac{\partial \mathbf{f}}{\partial \mathbf{x}} \Big|_{\mathbf{x}^{i-1}(s)} (\mathbf{x}^i(s) - \mathbf{x}^{i-1}(s)) \right) ds \quad (5.11)$$

Note that Eq. 5.11 is different from the original MCPI method in that the solution $\mathbf{x}^i(\tau)$ appears on both sides of the equation.

Assume that the n^{th} order Chebyshev expansions for the current iteration solution $\mathbf{x}^i(\tau)$, the previous iteration solution $\mathbf{x}^{i-1}(\tau)$, function \mathbf{f} , and the Jacobian matrix of \mathbf{f} , have the following forms

$$\mathbf{x}^i(\tau) = \sum_{k=0}^n \alpha_k T_k(\tau) \quad (5.12)$$

$$\mathbf{x}^{i-1}(\tau) = \sum_{k=0}^n \beta_k T_k(\tau) \quad (5.13)$$

$$\mathbf{f}(\mathbf{x}^{i-1}, \tau) = \sum_{k=0}^n \mathbf{F}_k T_k(\tau) \quad (5.14)$$

$$\frac{\partial \mathbf{f}}{\partial \mathbf{x}} \Big|_{\mathbf{x}^{i-1}} = \sum_{k=0}^n \mathbf{D}_k T_k(\tau) \quad (5.15)$$

Utilizing the multiplication property of the Chebyshev polynomials

$$T_m(\tau)T_n(\tau) = \frac{1}{2} (T_{m+n}(\tau) + T_{|m-n|}(\tau)) \quad (5.16)$$

for the second part of the integrand in Eq. 5.11, we may expect that a similar form of equations for solving for the Chebyshev coefficients can be obtained by using the same approach through which the original MCPI method is derived. However, com-

puting the Jacobian matrix and its Chebyshev approximations can be challenging. Furthermore, the linearization method couples the coefficient equations for different states, which can complicate the equations to solve for the coefficients tremendously, especially for high dimensional problems. We believe a general formula for solving for the coefficients is difficult to obtain if not impossible. More importantly, solving the resulting high dimensional equations will be very time-consuming.

We consider a special case which leads to a compact analytical equation for solving for the Chebyshev coefficients. The ordinary differential equation is a scalar case problem and all the Chebyshev coefficients D_k in Eq. 5.15 are treated as zero except the zero order term. Under these assumptions, Eq. 5.11 leads to

$$\begin{aligned}
& \frac{1}{2}\alpha_0 T_0(\tau) + \alpha_1 T_1(\tau) + \alpha_2 T_2(\tau) + \cdots + \alpha_N T_N(\tau) = \\
& \mathbf{x}(-1) + \int_{-1}^{\tau} \left(\frac{1}{2}(F_0 + \frac{1}{2}D_0(\alpha_0 - \beta_0))T_0(s) + \cdots + (F_N + \frac{1}{2}D_0(\alpha_N - \beta_N))T_N(s) \right) ds \\
& = \mathbf{x}(-1) + \frac{1}{2}(F_0 + \frac{1}{2}D_0(\alpha_0 - \beta_0))T_1(\tau) + \frac{1}{4}(F_1 + \frac{1}{2}D_0(\alpha_1 - \beta_1))T_2(\tau) + \cdots \\
& + \frac{1}{4}(F_N + \frac{1}{2}D_0(\alpha_N - \beta_N))\left(\frac{1}{N+1}T_{N+1}(\tau) - \frac{1}{N-1}T_{N-1}(\tau)\right) \tag{5.17}
\end{aligned}$$

Including the initial condition $x(-1) = x_0$ and equating the coefficients of the same order for both sides lead to the following linear equation for solving for the $(N + 1)$ Chebyshev coefficients

$$\begin{bmatrix} \frac{1}{2} & -1 & 1 & \cdots & (-1)^N \\ D_0 & -4 & -D_0 & \mathbf{0} & \mathbf{0} \\ 0 & D_0 & -8 & -D_0 & \mathbf{0} \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ \mathbf{0} & \mathbf{0} & \mathbf{0} & D_0 & -4N \end{bmatrix} \begin{bmatrix} \alpha_0 \\ \alpha_1 \\ \alpha_2 \\ \vdots \\ \alpha_N \end{bmatrix} = \begin{bmatrix} x_0 \\ 2(F_2 - F_0) + D_0(\beta_0 - \beta_2) \\ 2(F_3 - F_1) + D_0(\beta_1 - \beta_3) \\ \vdots \\ -F_{N-1} + D_0\beta_{N-1} \end{bmatrix} \tag{5.18}$$

In Eq. 5.18, $\mathbf{0}$ represents a zero vector with the appropriate dimension. Except for the

first row and the last row, the subset of the coefficient matrix has a special structure: the diagonal term of the i^{th} row is $-4(i-1)$, the entry before the diagonal term is D_0 , and the entry after the diagonal term is $-D_0$.

It is interesting to note that Eq. 5.18 is consistent with the equation derived by Norton [22]. Norton starts the derivation by using the following iteration form

$$\frac{dx^i}{d\tau} = f(x^i, \tau) + (x^i - x^{i-1}) \frac{\partial f}{\partial x} \Big|_{x^{i-1}} \quad (5.19)$$

Equation 5.18 can be generalized to solve BVP using the same assumptions for D_k . Both the zero and first order coefficients should be constrained by the boundary conditions.

D. Correction Control MCPI Methods

As shown in Chapter III, MCPI methods are essentially not stable when the domain of the independent variable is large. However, if we treat the sequence of the solutions generated from MCPI methods as a dynamic system response, it is possible that there exist some control algorithms that can manipulate the system behavior such that the new sequence of the solutions converges to the right solution even though the original sequence diverges.

Inspired by the proportional control, which is one of the most simple means of control, a proportional correction control MCPI method is designed as follows. Assume the original MCPI method generates a solution \mathbf{x}^- at the i^{th} step and generates a new solution \mathbf{x} at the $(i+1)^{\text{th}}$ step, and then the correction control MCPI method updates the solution at the $(i+1)^{\text{th}}$ step through

$$\mathbf{x}^+ = \mathbf{x}^- + C(\mathbf{x} - \mathbf{x}^-) \quad (5.20)$$

where C is a matrix that acts like the proportional gain. The matrix C can be constant or time varying, and can even be adaptively adjusted. Notice that when C becomes an identity matrix, Eq. 5.20 has the same form as the original MCPI method.

E. Numerical Examples

1. Characteristics of Convergence of MCPI Methods for Solving IVPs and BVPs

In Chapter III and Chapter IV, we focus on comparing the time efficiency and the accuracy of MCPI methods with other traditional methods. To illustrate the benefit of using the three techniques introduced in this chapter for improving the original MCPI methods, we first look at some convergent characteristics of MCPI methods when solving IVPs and BVPs.

The two-body problem discussed in Chapter III and the classical Lambert's problem presented in Chapter IV are the study cases. Again, the F and G solution provides a baseline for verifying the obtained results using MCPI methods. For the Lambert's problem case, we start by integrating the orbit for some time using some known initial position $\mathbf{r}(t_0)$ and velocity $\mathbf{v}(t_0)$ to obtain a position at the final time as $\mathbf{r}(t_f)$. Next we input the positions at the initial time and at the final time $\mathbf{r}(t_f)$ as the two-point boundary conditions, and utilize the MCPI method to solve for the required velocity at the initial time $\bar{\mathbf{v}}(t_0)$. The error of the MCPI method in this case is defined as $|\mathbf{e}| = |\bar{\mathbf{v}}(t_0) - \mathbf{v}(t_0)|$. For the case that the time of flight is 1/7 of the orbit, the MCPI method converges in 32 steps and the norm of the initial velocity error is $1.555 \times 10^{-13} km/s$.

Figures 59 and 63 show the relative errors at each iteration for the initial value problem and the two-point boundary value problem. With a zigzag fashion, the errors

are decreasing in a nearly exponential form for both cases. Figures 60 and 64 show the convergence rate of the errors with respect to the number of the iterations, where the convergence rate is defined as the ratio of the solution errors between two immediate iterations. Instead of having a linear form, which is the assumption for the Aitken process, the convergence rates are oscillating. Because the stopping criterion chosen for MCPI methods is based on the corrections between three immediate iterations, the correction histories and correction rates are also shown in Figs. 61, 62, 65, and 66. Similar to the error histories, the corrections between each iteration are decreasing almost exponentially, but the rates of the corrections have oscillating forms.

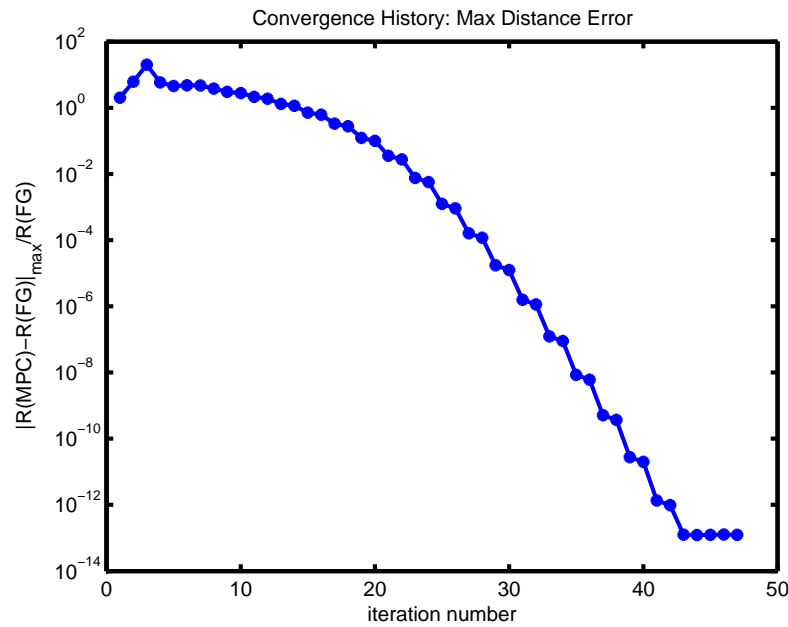


Fig. 59. Max relative distance error (IVP)

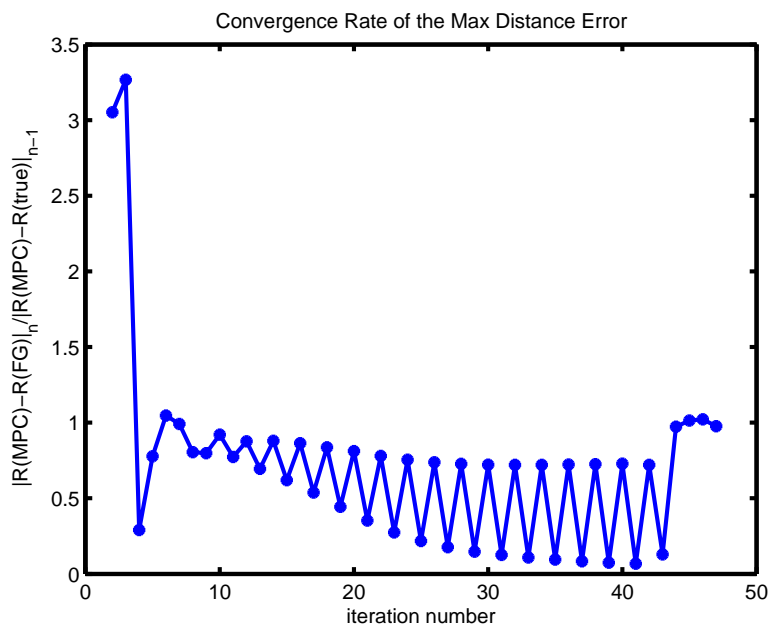


Fig. 60. Convergence rate of the max distance error (IVP)

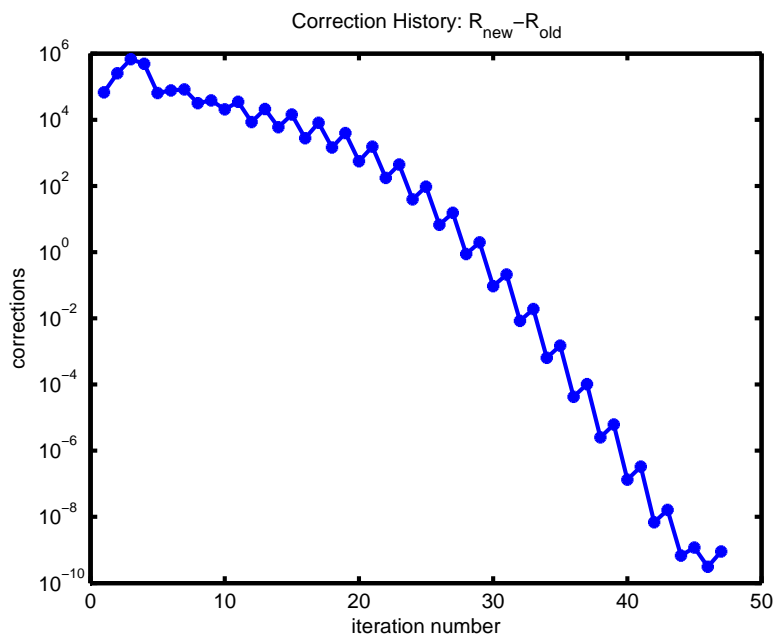


Fig. 61. Norm of the corrections (IVP)

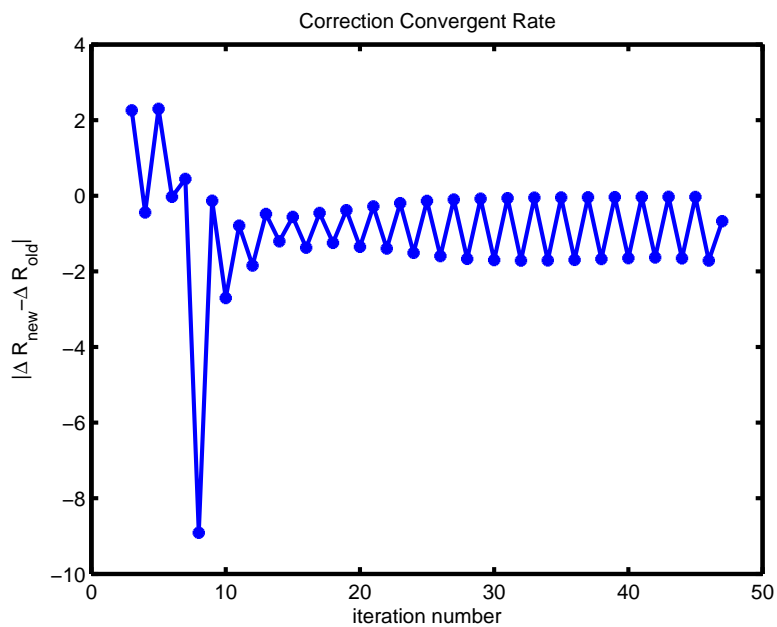


Fig. 62. Correction rate (IVP)

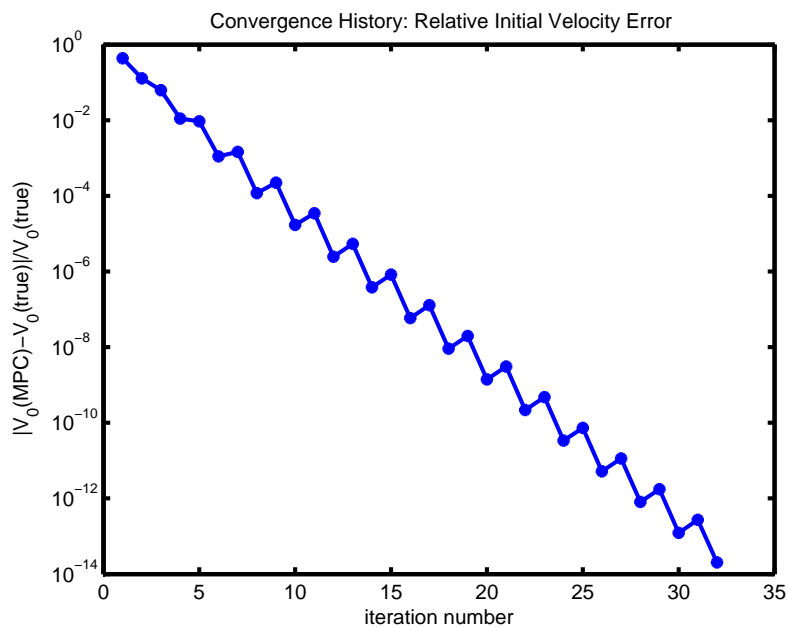


Fig. 63. Relative initial velocity error (BVP)

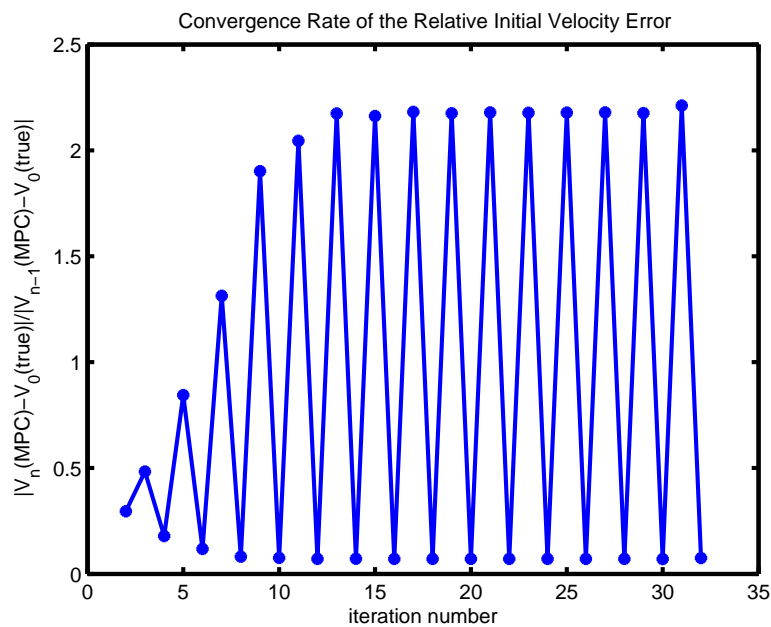


Fig. 64. Convergence rate of the initial velocity error (BVP)

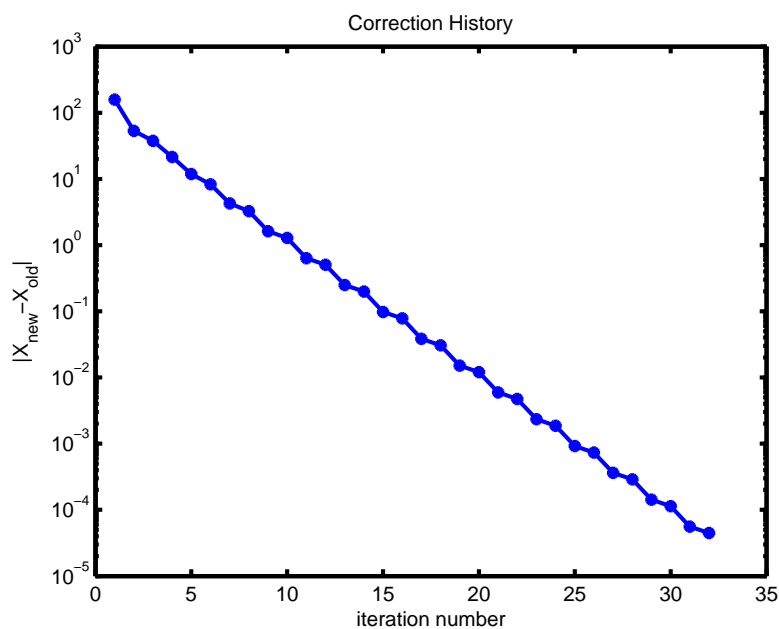


Fig. 65. Norm of the corrections (BVP)

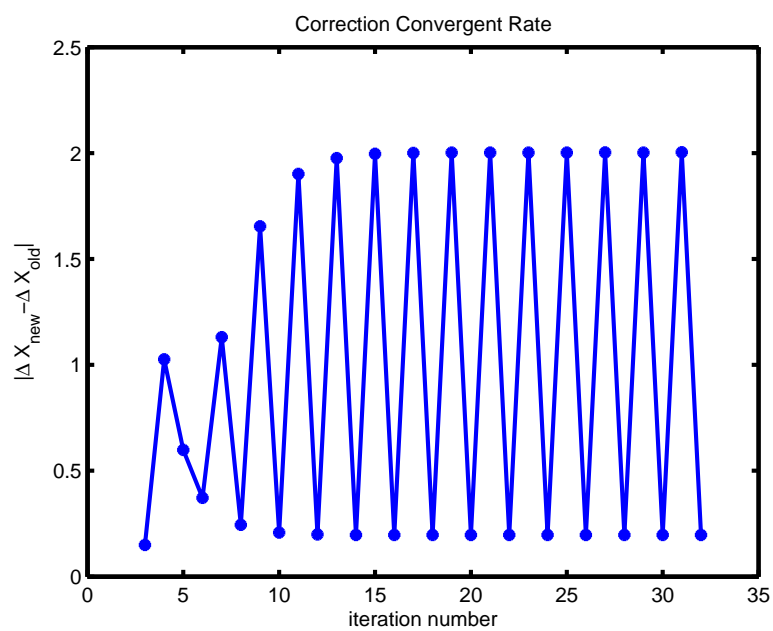


Fig. 66. Correction rate (BVP)

2. Using the Aitken MCPI Methods to Solve Lambert's Problem

Because Aitken's process is based on the assumption that the series converge in a geometric fashion, we should not expect that utilizing the Aitken's process for MCPI methods will always help to achieve significantly better performance. Although the Picard iteration itself has a geometric convergent behavior [31, 64], the proposed MCPI method has a much faster convergence rate than the linear form because Chebyshev polynomials are utilized. A rigorous and general transformation technique that is well-suited with MCPI methods, which have a hybrid nature of both the Picard iteration and Chebyshev polynomials, currently does not exist. But Aitken's process does appear to improve the convergence performance of MCPI methods.

We compare the performance of the original MCPI method and the Aitken MCPI method for Lambert's problem. The interval lengths are chosen from $\{0.1, 0.2, 0.3, 0.34, 0.35, 0.36, 0.37\}T_p$ where T_p is the orbital period. The required iteration numbers are shown in Fig. 67 and the CPU times are shown in Fig. 68. These plots show that, as the interval length increases, the benefit of using Aitken's process becomes more significant because it reduces both the number of iterations and CPU time. For the case where the interval length is $0.37T_p$, the CPU time of the MCPI method is two times the CPU time of the Aitken MCPI method and requires about three times of iterations. Additionally, when the time interval increases to $0.38T_p$ and $0.39T_p$, the original MCPI method diverges but the Aitken MCPI method converges after 282 and 432 iterations respectively.

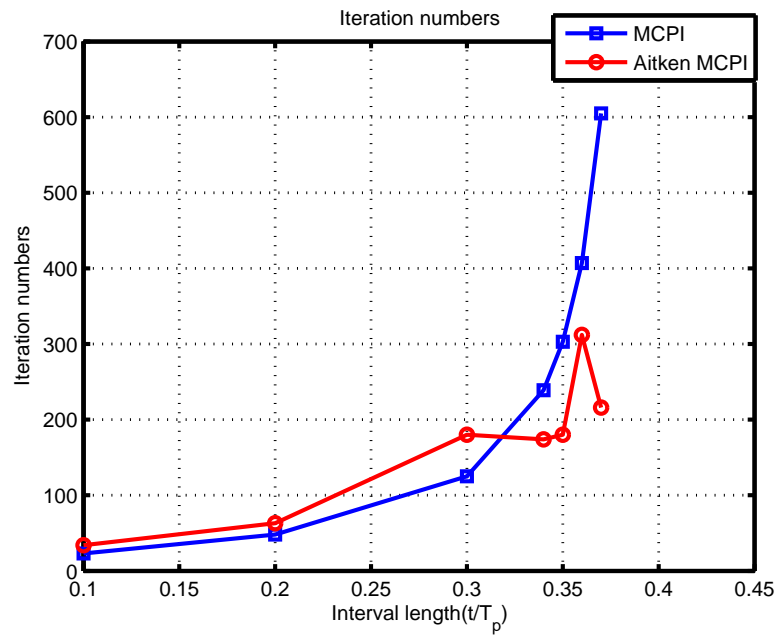


Fig. 67. Iteration number comparison (BVP)

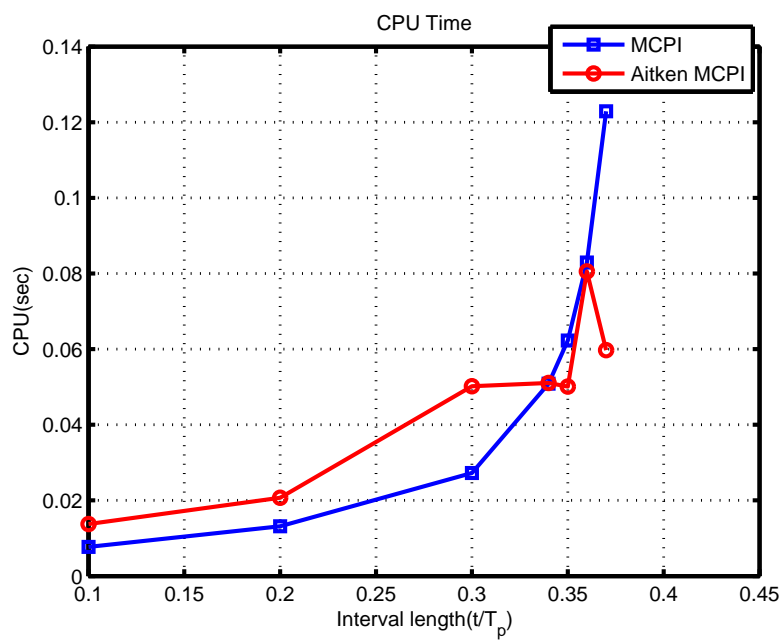


Fig. 68. CPU time comparison (BVP)

3. Using the Linearization MCPI Methods to Solve a Linear Problem

To satisfy the assumptions in Eq. 5.18, we consider an initial value problem

$$\frac{dx}{dt} = -\lambda x, t \in [0, T_f] \quad (5.21)$$

with an initial condition $x(0) = 1$. The analytical solution is $x(t) = e^{-\lambda t}$.

For the case when $T_f = 5$ seconds and $\lambda = 5$, the errors of the solutions using the MCPI method and the linearization MCPI method are shown in Figs. 69 and 70. The original MCPI method takes 333 iterations to find the solution in 0.02 seconds, whereas the linearization MCPI method takes 37 iterations to find the solution in 0.1 seconds. The order of the Chebyshev polynomial is chosen as 100 and the direct inversion approach is used to solve for the coefficients with Eq. 5.18. Although fewer iterations are required for the linearization MCPI method, solving the linear equations takes additional time, which causes the CPU time for the linearization MCPI method to be larger than the CPU time for the original MCPI method. However, increasing the final time T_f to ten seconds causes the original MCPI method to diverge, whereas the linearization MCPI method finds the solution using 58 iterations in 0.14 seconds. The solution errors are shown in Fig. 71.

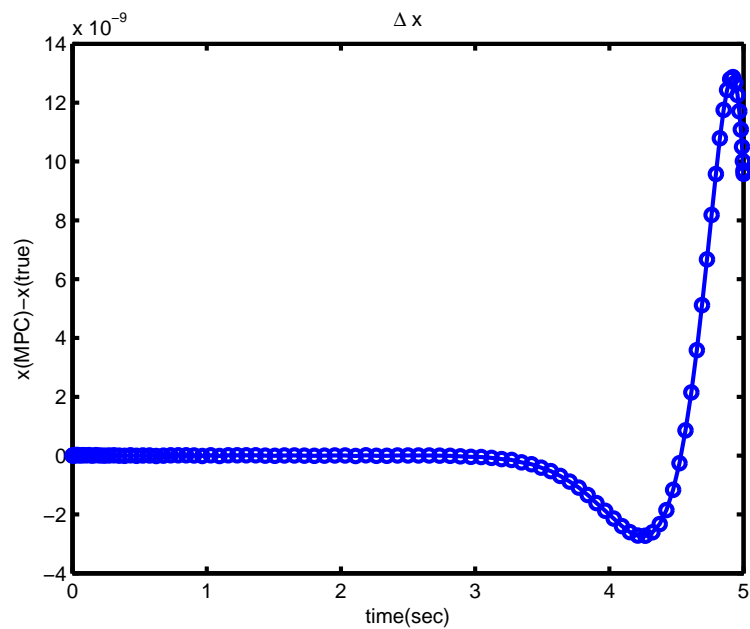


Fig. 69. Error history of the MCPI method ($T_f = 5s$)

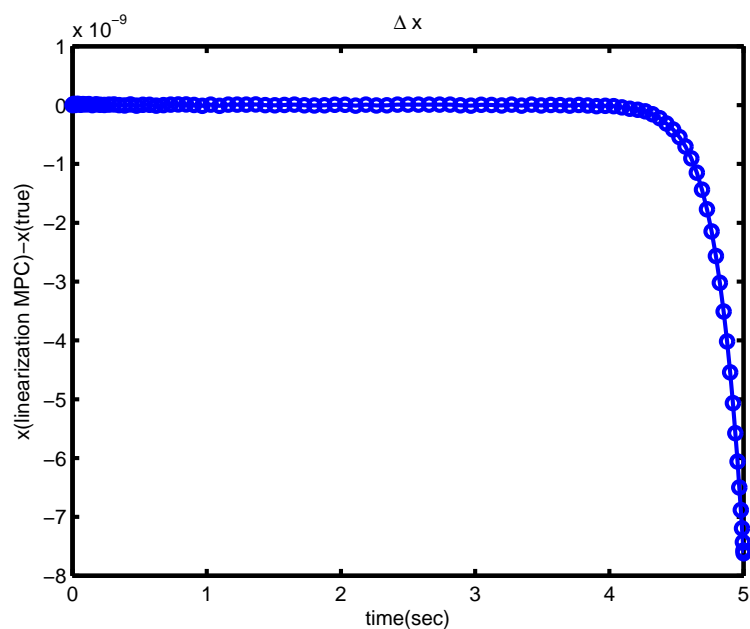


Fig. 70. Error history of the linearization MCPI method ($T_f = 5s$)

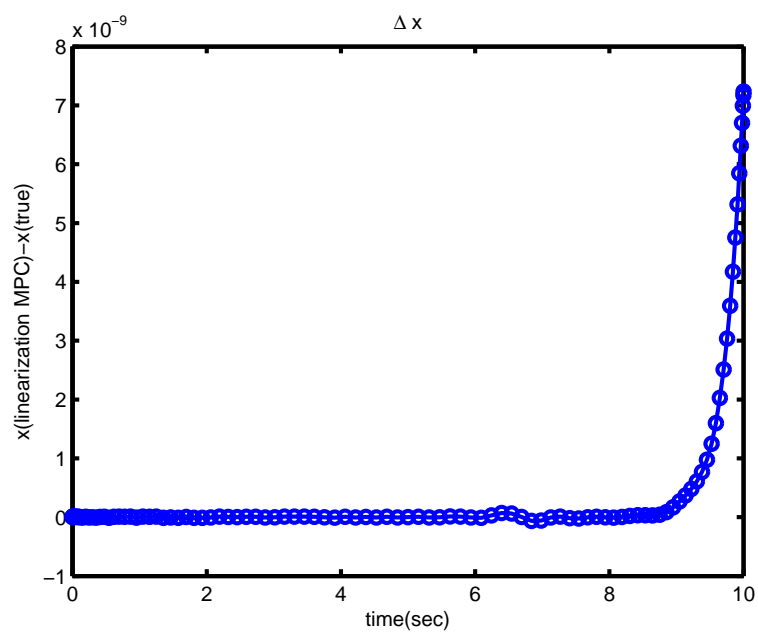


Fig. 71. Error history of the linearization MCPI method ($T_f = 10s$)

4. Using Correction Control MCPI Methods to Solve IVPs and BVPs

For the two-body problem discussed in Chapter III, the original MCPI method failed to converge after 1000 iterations using polynomials of order 400 and the stopping criterion of 10^{-6} . Figure 72 shows that the corrections of the sequence generated from the MCPI method are oscillating after about 200 iterations and the accuracy no longer improves. The relative position errors after the 1000 iterations are shown in Fig. 73. Choosing all the diagonal terms of the C matrix in Eq. 5.20 as 0.7, the correction control MCPI method converges after 240 iterations. This correction history is shown in Fig. 74, and the final relative position errors are shown in Fig. 75.

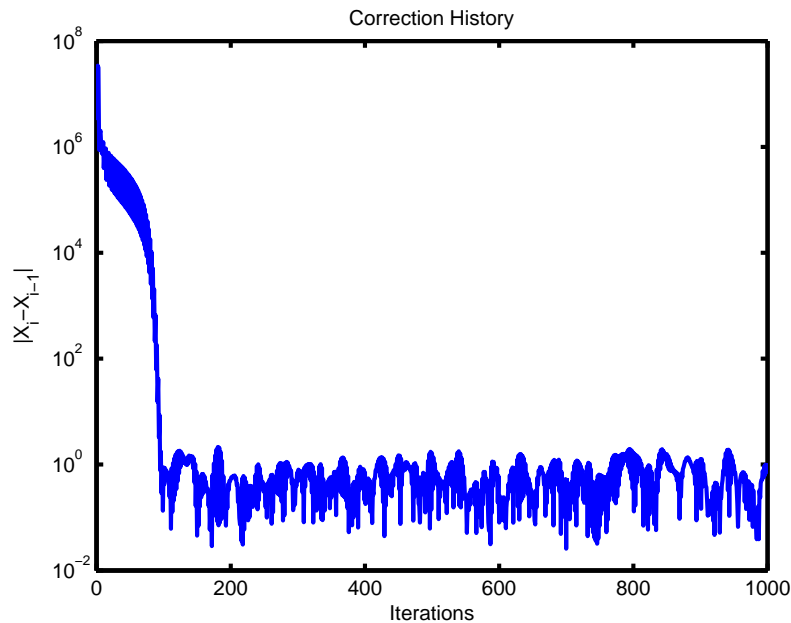


Fig. 72. Correction history using MCPI

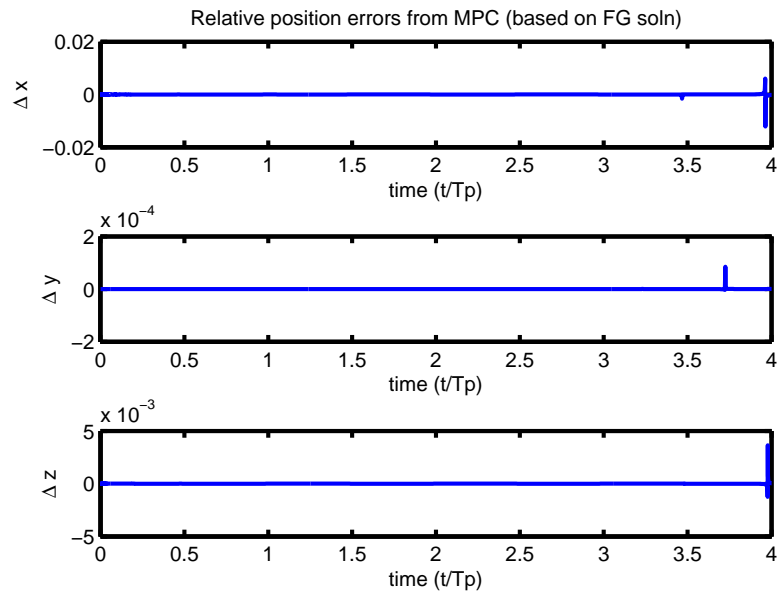


Fig. 73. Relative position errors using MCPI

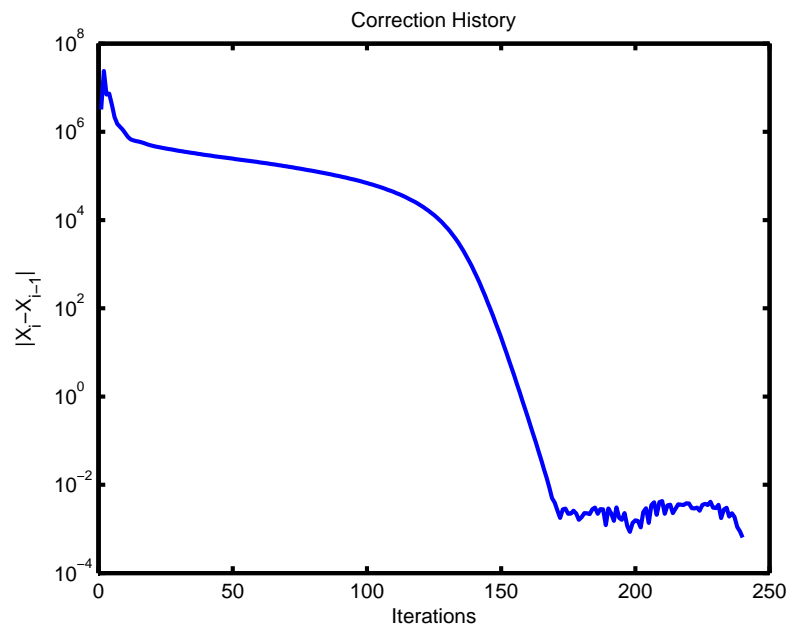


Fig. 74. Correction history using correction control MCPI

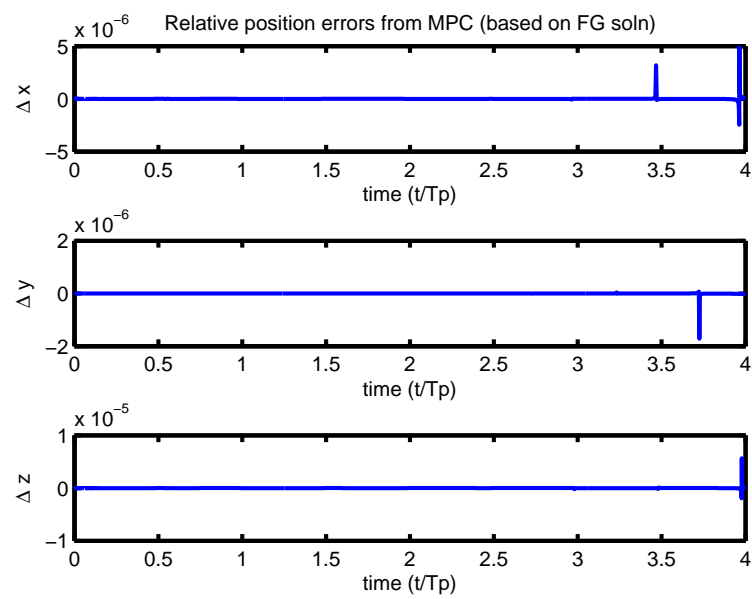


Fig. 75. Relative position errors using correction control MCPI

The Earth to Apophis optimal trajectory design problem is used here again to illustrate the benefit of using the correction control MCPI method to solve BVP. The original MCPI method only converges when the time of flight is no longer than 110 days. The correction control MCPI method expands this convergence domain significantly. Here the solution can be found for a time of flight of 150 days when the feedback gain C is chosen as a constant diagonal matrix with all the diagonal terms set to be 0.7. For this case, SNOPT converges in approximately five seconds with the boundary condition errors of magnitude 10^{-8} . The homotopy method finds the solution in seven and half seconds with the boundary condition errors of magnitude 10^{-13} . The MCPI method takes significantly less time to converge: 0.3 seconds in 226 iterations. Figures 76 to 81 show that the errors of the solutions for both the states and co-states are within the boundary of 10^{-9} .

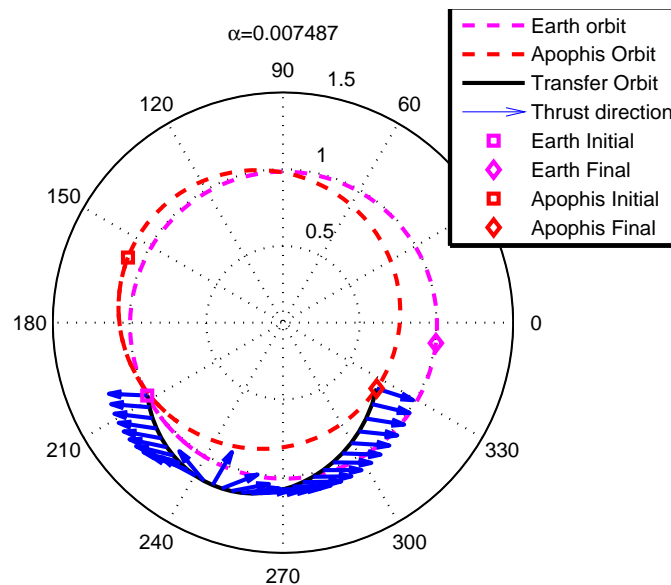


Fig. 76. Transfer orbit with time of flight 150 days

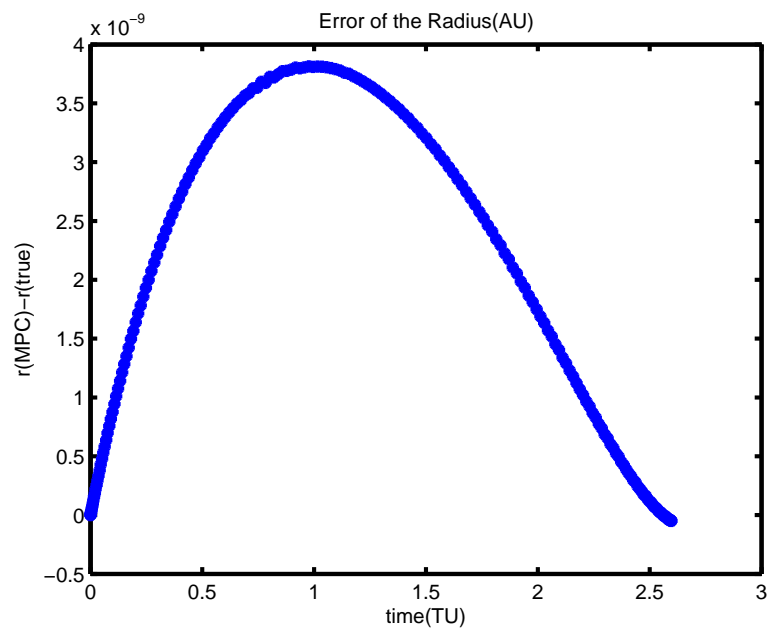


Fig. 77. Error of the radius (150 days)

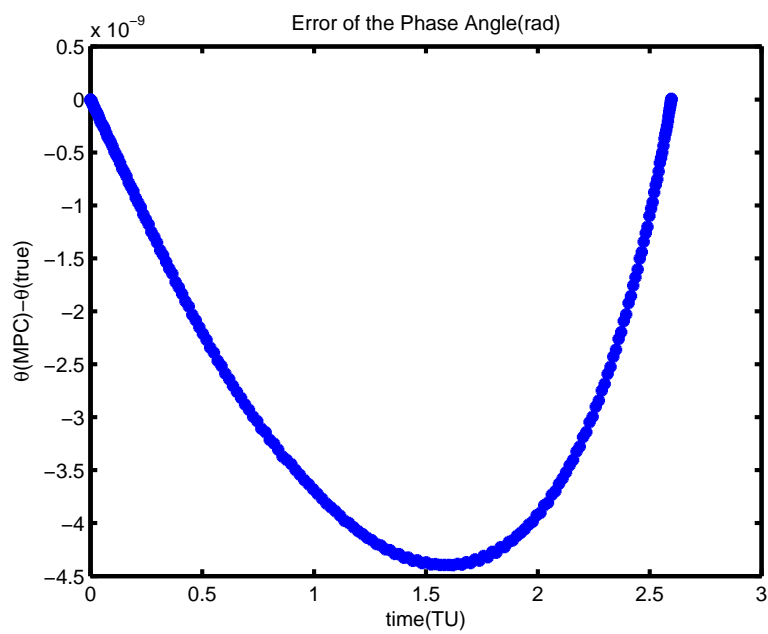


Fig. 78. Relative error of the phase angle (150 days)

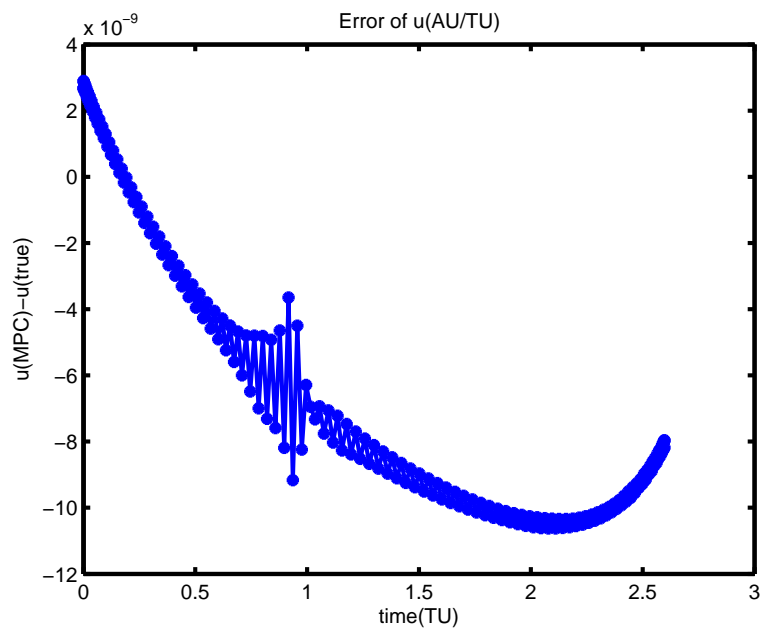


Fig. 79. Error of the radial velocity (150 days)

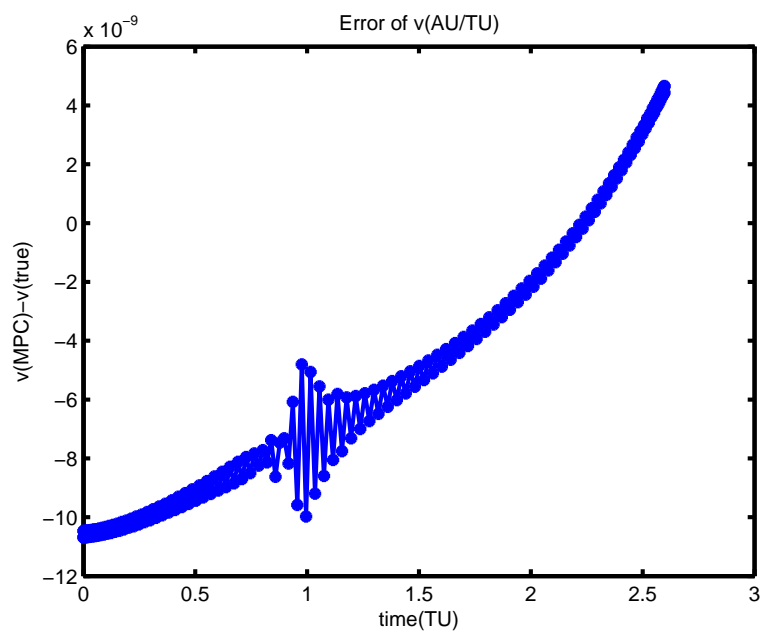


Fig. 80. Error of the tangential velocity (T150 days)

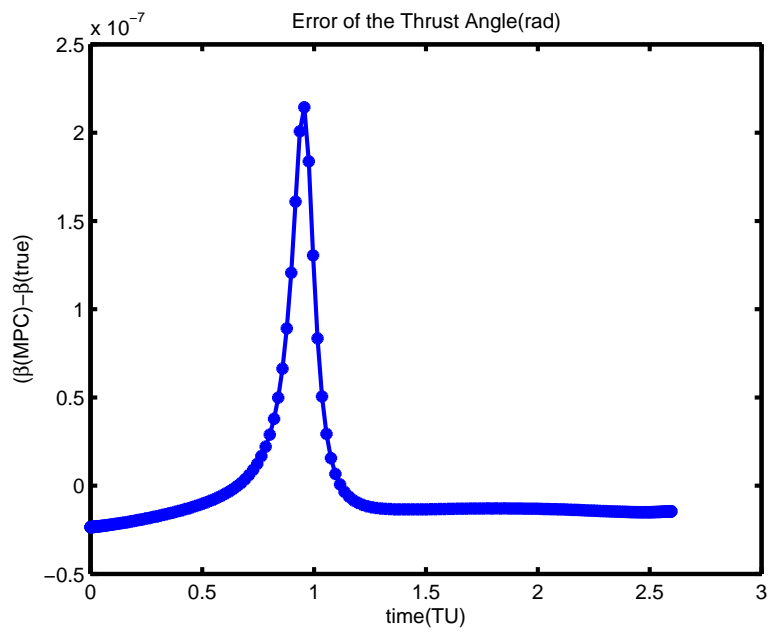


Fig. 81. Error of the thrust angle (150 days)

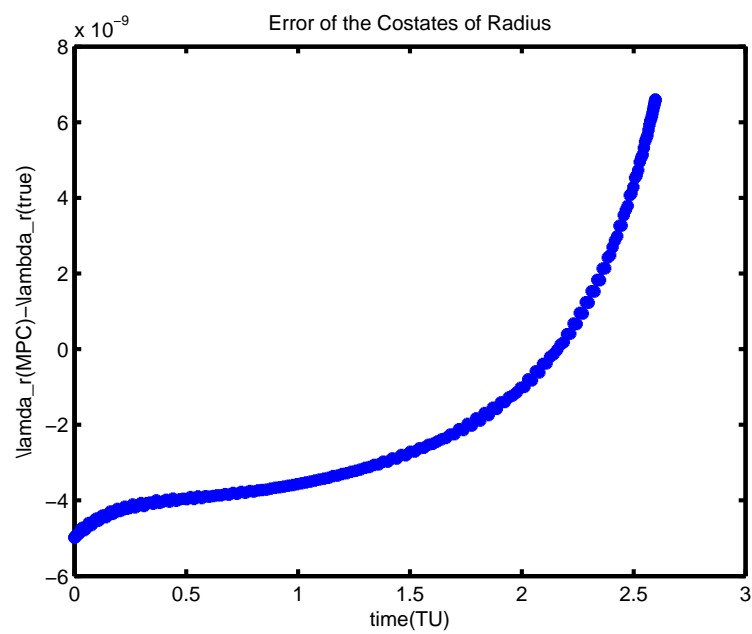


Fig. 82. Error of the costate about radius (150 days)

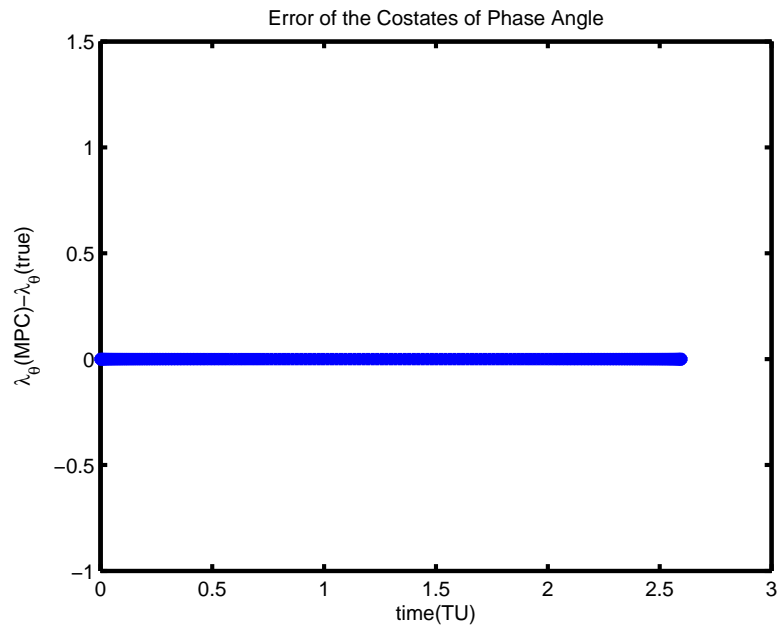


Fig. 83. Error of the costate about phase angle (150 days)

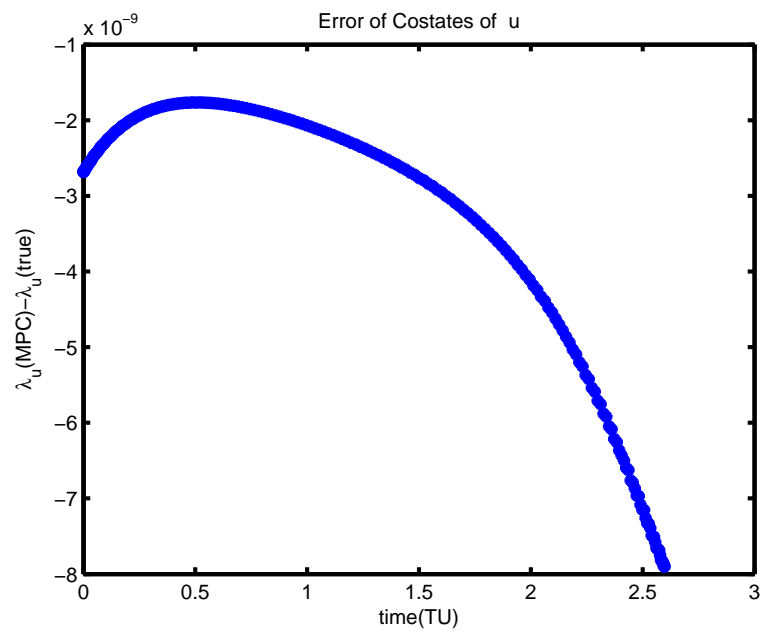


Fig. 84. Error of the costate about radial velocity (150 days)

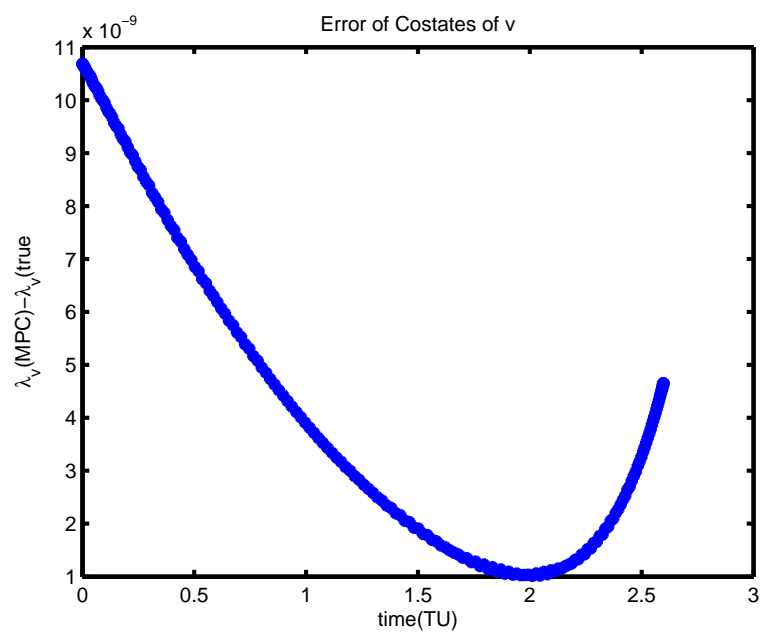


Fig. 85. Error of the costate about tangential velocity (150 days)

Similar to the importance of proportional gain in control of dynamic system behaviors, the gain matrix C is a key factor affecting the convergence behavior of the correction control MCPI method. Figure 86 shows the required iterations for different diagonal gain values, and Fig. 87 shows the CPU time for different chosen gains. These plots show that increasing the gain reduces both the number of iterations and the CPU time. However, too large of a gain will cause the methods to diverge again.

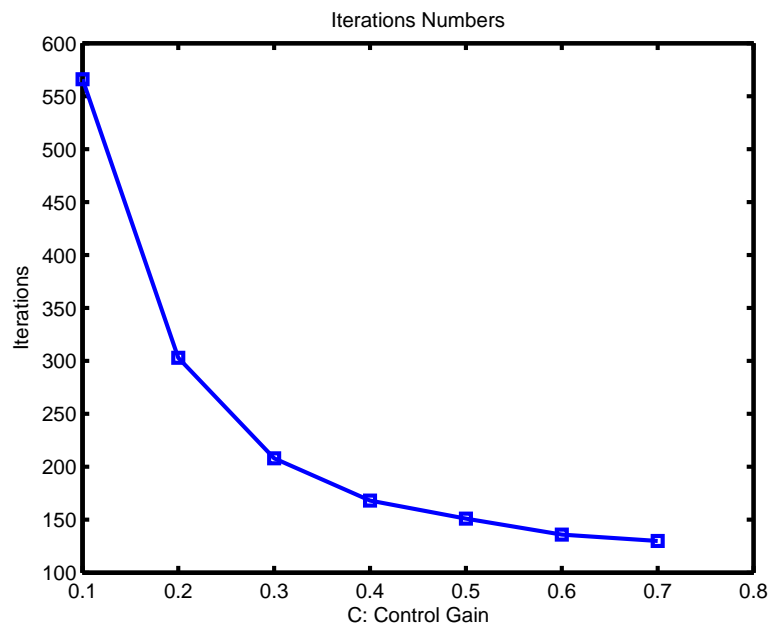


Fig. 86. Iteration number comparison (150 days)

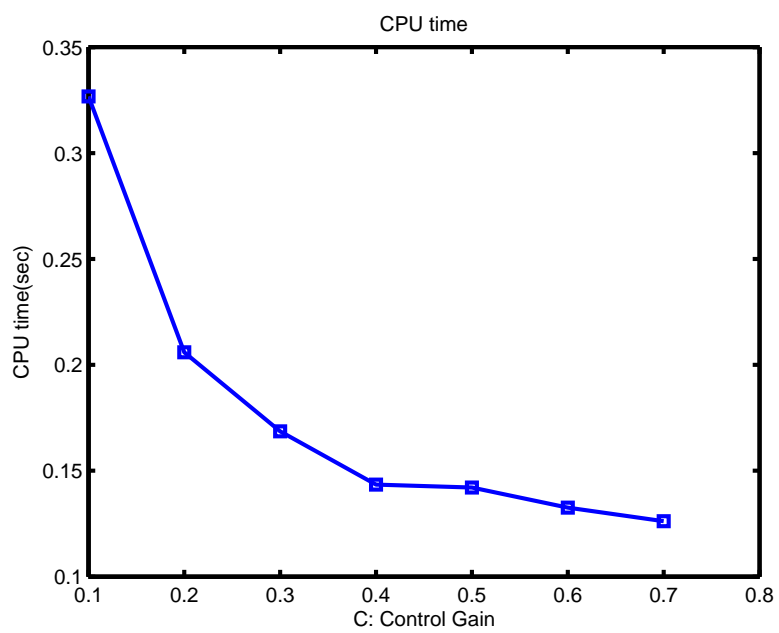


Fig. 87. CPU time comparison (150 days)

F. Summary

Although Aitken's process is designed for linearly convergent sequences, it is capable of improving MCPI methods, which do not converge in a geometrical fashion. The linearization approach can extend the convergence domain for the original MCPI methods, but it is difficult to apply to high-dimension nonlinear problems. The proportional correction control method is a very simple technique that can enlarge the convergence domain significantly. We suspect that more advanced control methodology can further enlarge the convergence domain. Such schemes may include choosing an optimal gain matrix that can vary during the iteration process as well as designing some more advanced control methods. Up through this chapter, all the numerical results have been obtained using the serial computation in MATLAB. The next chapter will discuss implementing MCPI methods on a graphics card, which provides a parallel computation environment for testing the proposed MCPI methods.

CHAPTER VI

IMPLEMENTATION OF MCPI METHODS USING GRAPHICS PROCESSING
UNITS

A. Introduction

In Chapter III and Chapter IV, MCPI methods were shown to achieve both improved accuracy and significant speedup when solving IVPs and BVPs in a serial computation environment. Due to the parallel structure of the methods, additional speedup is expected when the serial implementations are extended to a parallel computation environment. However, to achieve the most potential from using the parallel processors, the structure of the algorithm must be complementary to the architecture of the hardware. As warned by Thurber and Wald, any attempt to use a parallel computer with the wrong type of problems or methods is an exercise in futility [65]. Possible computer architectures can be classified into the following four categories according to the taxonomy introduced by Flynn [66].

- Single Instruction, Single Data stream (SISD): the computer architecture in which a single instruction is operated on a single stream of data. This type is basically the sequential computer where no parallelism is used.
- Single Instruction, Multiple Data streams (SIMD): the computer architecture in which a single instruction is operated on multiple streams of data simultaneously. One example of this type is the Graphics Processing Unit (GPU).
- Multiple Instruction, Single Data stream (MISD): the computer architecture in which multiple instructions are operated on single stream of data. This type of computer is mostly used for fault tolerance.

- Multiple Instruction, Multiple Data streams (MIMD): the computer architecture in which multiple instructions are operated on multiple streams of data simultaneously. One example of this type is the distributed system.

The proposed MCPI methods are well suited for large-scale data parallelism. Distributing the computation of both the force function evaluation at different time nodes and the coefficients of the Chebyshev polynomials for different orders to multiple processors is straightforward. Speedup is achieved when MCPI methods are implemented on SIMD type of parallel computers. For problems such as maintaining satellite catalog (see Chapter I) in which the force functions are highly computationally expensive and the dynamic systems require high order Chebyshev polynomials to approximate, significant speedups can be obtained.

Because of its rapid development and inexpensive price, the NVIDIA graphics card, one type of SIMD hardware, is chosen as the hardware for implementing MCPI methods in a parallel environment. This chapter first introduces GPU technology and discusses existing methods that use NVIDIA GPUs to solve IVPs. Next, the Compute Unified Device Architecture (CUDA) environment and CUBLAS (an implementation of the Basic Linear Algebra Subprograms (BLAS) on top of the NVIDIA CUDA)¹ are presented. The approach using CUDA and CUBLAS to implement MCPI methods is then presented. The performance of the graphic card first is tested with a matrix multiplication problem. We then present the results obtained using the GPU-accelerated MCPI methods to solve the small perturbation from the sinusoid motion problem and the zonal harmonic perturbation involved satellite motion propagation problem.

¹<http://developer.nvidia.com/object/gpucomputing.html> [retrieved May 7, 2010]

B. NVIDIA Graphics Processing Units

GPUs are traditionally specialized for three-dimensional graphics real-time rendering. Their highly parallel structure makes them more effective than the general-purpose CPU for parallel computation. With the recent programmability enhancements associated with some development environment such as NVIDIA CUDA, GPUs have helped to solve complicated scientific computation problems in various fields such as molecular modeling, magnetic resonance imaging, N-body simulations, and evolutionary algorithms [67, 68, 69, 70]. NVIDIA GPUs represent the state of the art of General Purpose GPUs (GPGPU). The NVIDIA GPU consists of a set of multiprocessors that can be used by the graphic cards to solve many complicated computational problems in addition to the traditional 3D graphics functions.

One important reason for the successful application of GPUs to these problems is the amazing rate at which the GPU technology is developing, dwarfing even Moore's law that applies broadly to advances in computer hardware and computing speed. GPU-accelerated computational power nearly doubles every six months while CPU computational power doubles about every 18 months [71]. GPUs are relatively inexpensive; their competitive price is mainly due to mass production related to the high demand of the commercial gaming industry. As a consequence, we have entered the age of personal supercomputing, and it is possible to run large complicated problems on a personal computer. Previously this was impractical due to the lengthy processing time required. For example, the current GeForce GTX 280 GPU achieves $141.7GB/s$ memory bandwidth and 933 Giga Floating point Operations Per Second (GFLOPS) for single-precision computation. Thus, a 0.93 Terraflop personal computer is now available through the addition of a graphics card costing only a few hundred dollars.

Most current GPU-based applications for solving IVPs use simple forward in-

tegration algorithms for dynamic system simulations such as the first-order Euler integration method or the second-order leapfrog-Verlet integrator [72]. However, the accuracy of these low-order integration methods is often low. Therefore we must extend beyond these elementary applications and consider more powerful differential equation solvers. Additionally, both Euler and leapfrog-Verlet methods are essentially sequential, not parallel algorithms. Although when solving an all-pairs N-body problem by using the GPU techniques [72], Nyland et al. have achieved 50 times as fast as a highly tuned serial implementation or 250 times faster than a portable C implementation, it is safe to predict that an efficient parallelized integration algorithm such as the proposed MCPI methods can further improve the overall performance.

C. CUDA Environment and CUBLAS Toolbox

CUDA is a general purpose parallel computing architecture developed by NVIDIA². It is the computing engine that allows developers to use NVIDIA GPUs to conduct high performance computations. The programming mode of CUDA includes kernels, blocks, and threads. A kernel is executed as a grid of blocks, and a block is a batch of threads that can cooperate with each other by sharing data through shared memory. Threads belonging to one block can execute synchronically while threads belonging to different blocks can not cooperate. Since NVIDIA GPUs are based on SIMD architecture, CUDA achieves its best performance when a single operation is running simultaneously on many threads, which is well-suited for MCPI methods. Compared with CPU threads, CUDA threads are extremely lightweight with little creation overhead and can be switched instantly. Furthermore, while multi-core CPUs usually can have only a few threads, the current maximum number of threads for one

²<http://developer.nvidia.com/object/gpucomputing.html> [retrieved May 7, 2010]

grid on a graphics card is tens of thousands.

The CUDA environment provides two levels of application programming interface (API) functions: a low-level CUDA driver API and a high-level CUDA runtime API. These APIs provide five categories of functions for users to use: general functions such as the clock function and the standard math functions, device functions, memory management functions, flow control functions, and interface functions with OpenGL and Direct3D.

CUDA code is compiled by the NVIDIA C Compiler (NVCC). The code that runs on the CPU is called the host code and the code that runs on the GPU is called the device code. The NVCC separates the host code to its own compiler and links it with the compiled device code during the run time. To achieve the most potential using CUDA, in addition to a good knowledge of the general parallel computation theory, the developers need to understand some specific characteristics of the CUDA performance. For examples, the developers should take advantage of the much faster shared memory rather than the global memory, use memory coalescing to reduce latency, and optimize the program to avoid memory bank conflicts and divergent branching issues.

Three levels of BLAS functions are included in CUBLAS. Level one is for scalar, vector, and vector-vector operations; level two is for vector-matrix operations; and level three is for matrix-matrix operations. The CUBLAS toolbox includes both single precision functions and double precision functions. Users should make sure that the computing capability of their graphic cards is at least 1.3 in order to use the double precision functions correctly. Notice that the functions in CUBLAS are not always the fastest realization of the algorithms because they are written for many possible cases. For example, the matrix multiplication function *cublasSgemm* considers additional cases where one or two of the input matrices are transposed, the multiplication result

is multiplied by a scalar, and another summation is involved. Still we choose to use CUBLAS because the implemented functions have been tested on different graphic cards and are robust to various programming errors.

The way we implement MCPI method is shown in Fig. 88, where *GPU* denotes the code is running on the device (graphics card), *CPU* denotes the code is running on the host, and *CUBLAS* denotes the CUBLAS toolbox is used to do the computation. Notice that the coefficient matrices C_x and C_α are constant once the order of polynomials is defined. They are calculated before the iteration starts.

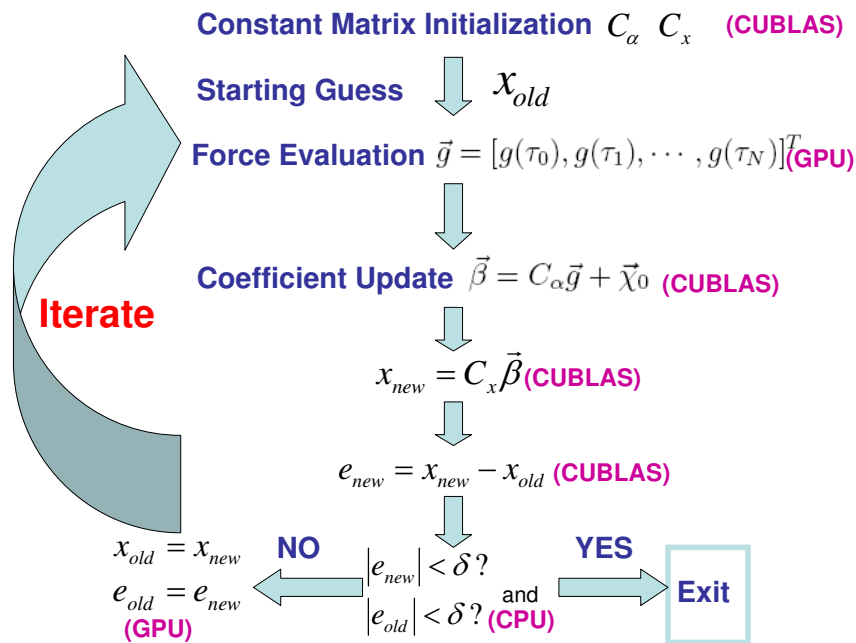


Fig. 88. Implementation of MCPI methods using CUDA and CUBLAS

D. Graphic Card Testing

The settings of the computer and the development environment used are the following

- Intel(R) Pentium(R) D CPU 3.4GHz, 3.4GHz, 2.0GB of RAM
- Windows XP Operating System
- NVIDIA GeForce 9400GT Graphics Card
- MATLAB R2009b
- Microsoft Visual Studio 2005

The theoretical memory bandwidth for GeForce 9400GT is 12.8 GB/s and the peak GFLOPS is 67.2 for single precision computation. The runtime example of a single precision matrix multiplication indicates that the bandwidth is 7.9419GB/s, and the maximum GFLOPS is 25.

Due to the vector-matrix characteristic form of MCPI methods, we compare the performance of three different computation environments for matrix multiplication first. The results are shown in Table I, where

- N: dimension of the square matrices
- MB(A+B+C): size of A+B+C in Megabytes
- MT: MATLAB computation time
- h2d: data transfer time from the host to the device
- kernel: computation time using cublasSgemm function
- d2h: data transfer time from the device to the host
- $\frac{serial-C}{GPU}$: ratio of the computation time between that of the serial C code and cublasSgemm

- $\frac{MT}{GPU}$: ratio of the computation time between that of the MATLAB code and cublasSgemm
- GFLOPS is calculated from

$$GFLOPS = \frac{2N^3}{time(kernel)} 10^6$$

Because the MATLAB clock function can not return the CPU time when the computation time is less than $1ms$, “N/A” is shown for these cases in the table.

Table I. Time comparison of matrix multiplication $A \times B = C$

N	MB(A+B+C)	CPU(ms)	MT(ms)	GPU(ms)				$\frac{serial-C}{GPU}$	$\frac{MT}{GPU}$
				h2d	Kernel	d2h	GFLOPs		
256	2	141	N/A	0.5	2	0.3	20	70	N/A
512	4	1359	47	2	11	1	24.4	124	4.3
1000	11.4	11266	297	5	184	3	11	62	1.6
1024	12	35594	313	6	85	3	25.2	419	3.7

From Table I, we conclude that

- For a matrix dimension that is a multiple of 32, the achieved GFLOPS of the GPU are above 20; for matrix dimension that is not a multiple of 32, the achieved GFLOPS of GPU is relatively low. This is because CUDA executes an operation as half warp, which has 16 threads. The device code achieves its greatest efficiency when the number of the threads per block is a multiple of warp size.
- The larger the dimension of the matrix, the greater the speedups are achieved from using GPUs, as compared with both serial C code and MATLAB code.

The speedups of using GPU over MATLAB are three to four for the matrix dimension that is a multiple of 32; the speedups of using GPU over serial C code are in the range of 60 to 400.

E. GPU-Accelerated MCPI Methods for Solving the Small Perturbation from the Sinusoid Motion Problem

1. Sequential MCPI and GPU-accelerated MCPI in C

Here we choose the small perturbation from the sinusoid motion discussed in Chapter III as a case study. The results are shown in Table II, where

- N: order of the polynomials
- C_x, C_α : computation time for the two constant matrices
- Iteration: computation time during the iterations
- h2d: matrix transferring time from the host to the device
- d2h: matrix transferring time from the device to the host
- Overall: total computation time: h2d+(C_x, C_α)+ iteration+d2h

Table II. Time comparison of sequential MCPI and GPU-accelerated MCPI in C

N	CPU time(ms)			GPU time(ms)					$\frac{CPU}{GPU}$
	C_x, C_α	Iteration	Overall	h2d	C_x, C_α	Iteration	d2h	Overall	
511	5391	16	5407	4	43	4	0.3	51	106
1000	57766	125	57891	11	747	29	0.3	787.3	73.5
1023	163797	156	163953	12	340	8	0.3	360.3	455

Table II shows that

- For the largest polynomials of order 1023, which are beneficial to CUDA computation, the speedups achieved by using the GPU are more than four hundred.
- The matrix multiplication operation is the dominate cause for these speedups.
- The speedup from the iteration part is about four for $N = 511$ and $N = 1000$; for $N = 1023$ the gain is about twenty.

2. GPU-Accelerated MCPI in C and Sequential MCPI in MATLAB

Because the vector and matrix operations using the CPU are developed “in-house”, they may be significantly less efficient than the BLAS functions. Thus, the previous comparison is biased to the GPU implementation. Because the vector and matrix operations in MATLAB are highly efficient, we choose to compare the implementation of MCPI methods between using GPU in C environment and using sequential MATLAB with single precision calculation. Table III shows that the speedups achieved by using the GPU card over the MATLAB implementation are in the range of 2 to 4.5. From Table III, some important conclusions can be drawn, such as

Table III. Time comparison of sequential MCPI in MATLAB and GPU-accelerated MCPI in C

N	MATLAB time(ms)			GPU time(ms)					$\frac{MT}{GPU}$
	C_x, C_α	Iteration	Overall	h2d	C_x, C_α	Iteration	d2h	Overall	
511	219	N/A	219	4	43	4	0.3	51	2
1000	1531	51	1582	11	747	29	0.3	787.3	2.0
1023	1578	31	1609	12	340	8	0.3	360.3	4.5

- The higher the order of polynomials, the greater the gain can be obtained from using the GPU implementation.
- For the matrix size which lends itself to GPU computation, the speedup accomplished by using GPU is even higher. When the order of the polynomial is chosen to be 1023, which leads to the matrix size of 1024, the greatest speedup results.

F. GPU-accelerated MCPI Methods for Zonal Harmonic Perturbation Involved Satellite Motion Propagation Problems

The zonal harmonic perturbation satellite motion propagation problem discussed in Chapter III is chosen as another study case. The simulation results are shown in Figs. 89 to 93. Several important observations are

- As can be seen from Figs. 89 and 90, compared with the sequential code in MATLAB, the computation time required by the GPU-accelerated MCPI code in C does not change significantly when higher order zonal harmonic perturbations are included.
- Figures 91 and 92 show that the speedups obtained from the GPU-accelerated MCPI code increase as either the perturbation forces become more complicated or higher order polynomials are used.
- Figure 93 shows the relative differences of the solutions between these two implementations are below 10^{-6} . Considering the graphic card we use is only capable of single precision computation, we expect the two obtained solutions will match better once we implement the GPU-accelerated MCPI code using an advanced card that can conduct double precision computation.

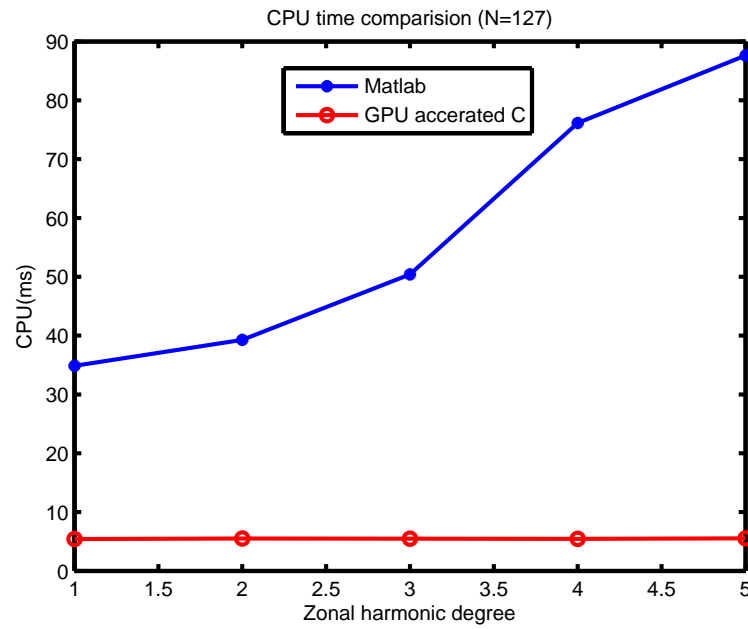


Fig. 89. Computation time of GPU-accelerated MCPI and MATLAB MCPI (N=127)

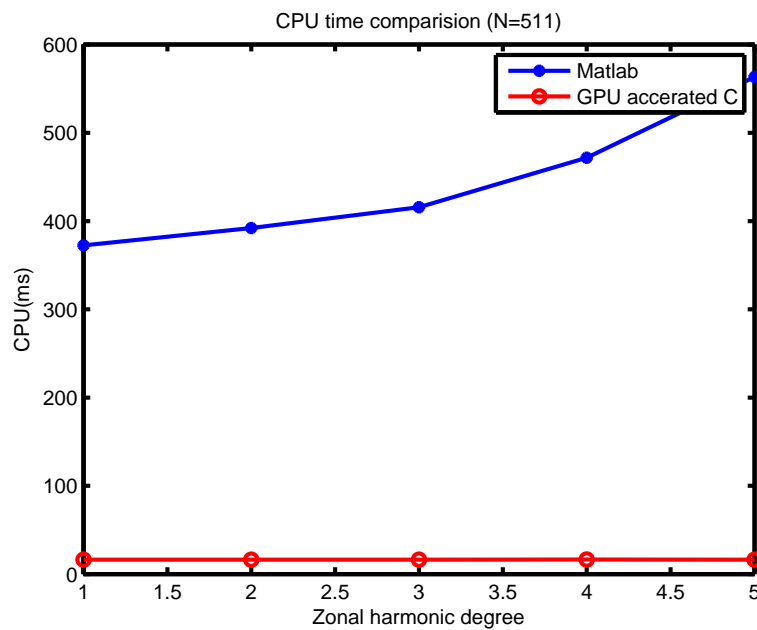


Fig. 90. Computation time of GPU-accelerated MCPI and MATLAB MCPI (N=511)

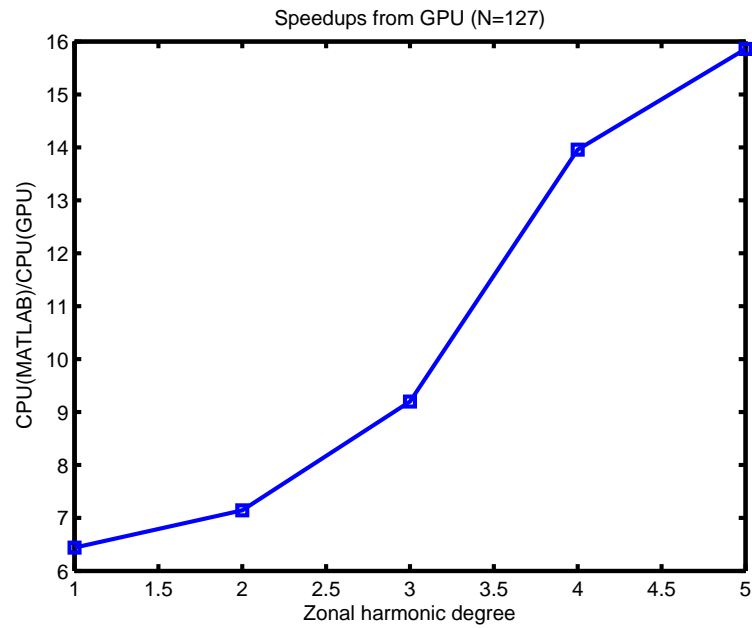


Fig. 91. Speedup of GPU-accelerated MCPI over MATLAB MCPI (N=127)

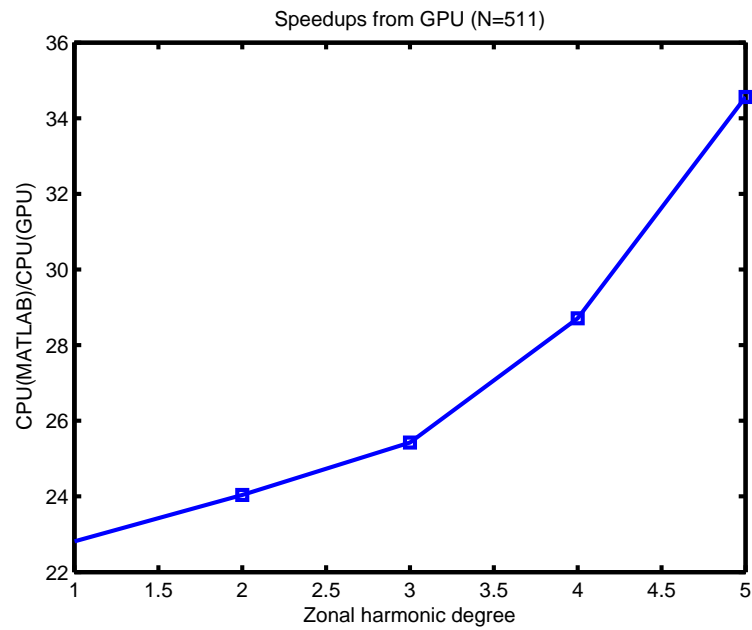


Fig. 92. Speedup of GPU-accelerated MCPI over MATLAB MCPI (N=511)

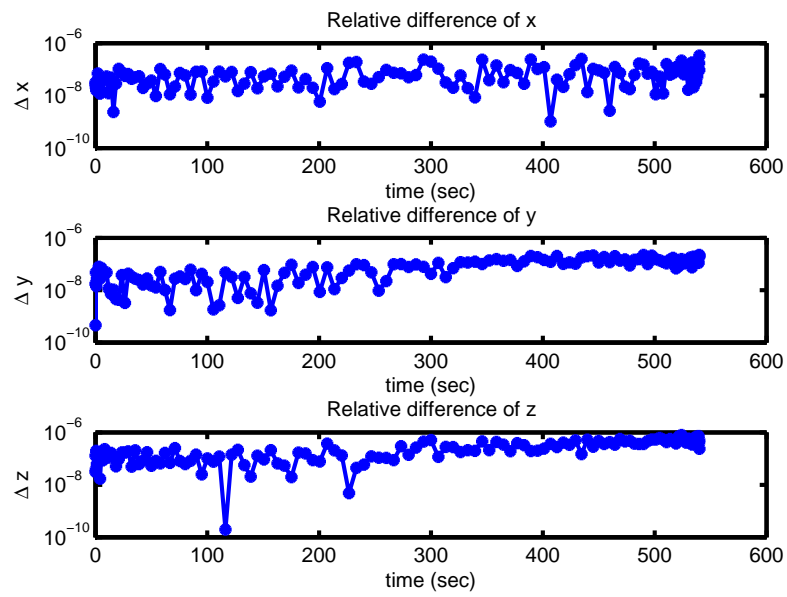


Fig. 93. Relative position difference between GPU-accelerated MCPI and MATLAB MCPI (N=511)

G. Summary

The recently developed GPU technology provides a SIMD parallel architecture for implementing MCPI methods. Although how much speedup can be achieved from the parallel implementation will depend on the problems themselves, this chapter has showed the power of parallel implementation of serial MCPI methods. For the small perturbation from the sinusoid motion problem where the force function is relatively simple, the speedups gained by using GPUs over the reference MATLAB solutions are in the range of two to five. For the satellite motion propagation problems, the more complicated the force functions and the higher order polynomials used, the larger the speedups that are achieved when using GPUs. The largest speedup for this study is 34.6, which is obtained when we use 511th order Chebyshev polynomials and include up to 5th order zonal harmonic perturbations.

CHAPTER VII

CONCLUSIONS

This dissertation has presented a unified framework for applying MCPI methods to solve IVPs and BVPs. The accuracy, efficiency, and parallel structure of MCPI methods are illustrated through several important celestial mechanics problems.

For IVPs, MCPI methods solve problems over arbitrary finite time intervals using a piecewise approach. For the small perturbation from the sinusoid motion and the satellite motion propagation problems, MCPI methods achieve better accuracy and speedups when compared with ODE45 in MATLAB, even prior to parallel implementation. Larger speedups are obtained when MCPI methods are implemented using GPU technology. The following two techniques have been shown to further speed up MCPI methods:

1. To solve IVPs with systems that are described by higher-order differential equations, MCPI methods do not require transforming the original equations to their first-order forms. MCPI methods can solve for the lower-order states directly without the need to integrate the higher-order information during iterations. The higher-order information can be obtained conveniently afterwards by using the properties of Chebyshev polynomials.
2. For IVPs that can be treated as perturbed motion from a reference trajectory, using MCPI methods to solve for the perturbation motion reduces the computation time while maintaining the solution accuracy.

For BVPs, MCPI methods find solutions that satisfy both dynamic equation constraints and boundary conditions with high accuracy. Compared with several other BVP solvers, MCPI methods are computationally efficient and not initial guess sen-

sitive. Using MCPI methods to solve optimal control problems through Pontryagin's principle, optimality and accuracy are guaranteed. For the cases where MCPI methods diverge, the proposed nonlinear transformations, the linearization approach, and correction control methods help to enlarge the convergence domain.

As one motivation for this study is the space catalog problem, we summarize the following results obtained when MCPI methods are used to solve the satellite motion propagation problems:

1. MCPI methods can integrate the satellite motion for arbitrarily long times. When compared with ODE45 in MATLAB, the longer the integration time, the larger the speedups achieved by MCPI methods.
2. To integrate the satellite motion for sixteen revolutions with only the inverse-square gravity field, using MCPI methods, one achieves up to three orders of magnitude better accuracy with more than one magnitude of speedups over ODE45 prior to parallel implementation. MCPI methods that only integrate position states, obtaining the velocity information afterwards, achieve up to two orders of magnitude speedups over ODE45.
3. Although the speedup of the sequential realization of MCPI methods drops when the degree of the included zonal harmonic perturbation forces increase, a GPU implementation results in significant speedups.

Considering that both sufficient accuracy and significant speedups are accomplished by the proposed MCPI methods prior to parallel implementation, and that parallel implementation achieves larger speedups when the force functions become more complicated and higher order polynomials are used to approximate the solutions, we claim that utilizing MCPI methods results in both accuracy and efficiency for solving satellite motion propagation problems.

With the advantages of MCPI methods in solving IVPs and BVPs, as shown in the dissertation, the proposed MCPI methods, especially given their inherently parallel structure, will contribute significantly to dynamic system analysis and controlling. Some suggestions for the future study include the following three fields:

1. Although the current study has shown that a parallel implementation of MCPI methods brings speedups over forward integration methods, further study could focus on specific applications and likely achieve even greater speedups than what we have obtained in this dissertation. For the space catalog mission, integrating the motion for 150,000 satellites requires another level of parallelism. Advanced graphic cards and other more powerful parallel computation environments such as clusters or distributed computing should be studied. Although this dissertation has provided some insight on how to choose the segment step size, polynomial order, and stopping criterion, these issues should be further studied for specific applications to achieve the potential of the MCPI methodology.
2. Through the optimal trajectory design problem, we have shown that once MCPI methods converge, they solve the problem more than one magnitude faster than several other popular optimizers prior parallel implementation. Considering the importance of integrating the dynamic systems when solving optimal control problems with Pontryagin's principle and the fact that MCPI methods are very efficient in solving IVPs, the superior performance of MCPI methods is not surprising. More efficiency is obtained because MCPI methods do not require gradient information for solving BVPs. The potential of MCPI methods for solving optimal control problems with the inherently parallel structure is huge. In addition to using MCPI methods to solve other types of optimal control problems, further study could focus on developing more powerful strategies for

enlarging the convergence domain.

3. It is through utilizing Chebyshev polynomials that the traditional Picard iterations converge fast, compute efficiently, and acquire a parallel structure. However, these characteristics could also be obtained by using other orthogonal basis functions. Developing a general strategy where Picard iterations are combined with other basis functions is valuable, as different basis functions may have different convergence properties. It is possible that for some specific problems, using one set of basis functions makes the Picard iteration diverge while using another set of basis functions leads to a converged Picard iteration. Furthermore, as has been shown, choosing a good reference trajectory and integrating only the perturbation motion will speed up the original MCPI methods. Similarly, choosing an appropriate basis function for a specific application will also contribute to greater speedups.

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APPENDIX A

PICARD ITERATION

For a scalar case first order differential equation

$$\frac{dx}{dt} = f(t, x) \quad (\text{A.1})$$

with an initial condition

$$x(t_0) = x_0 \quad (\text{A.2})$$

Picard iteration generates a chain of solutions $x^i(t)$ ($i = 1, 2, \dots, \infty$) by using

$$x^i(t) = x(t_0) + \int_{t_0}^t f(x^{i-1}, s) ds \quad (\text{A.3})$$

In a domain D surrounding the point (t_0, x_0) and defined by the inequalities

$$|t - t_0| \leq a, |x - x_0| \leq b \quad (\text{A.4})$$

if $f(t, x)$ is a one-valued continuous function of x and t and if $f(t, x)$ satisfies the Lipschitz condition, then the sequence in Eq. A.3 does indeed converge to a unique and continuous solution that satisfies the differential Eq. A.1. Lipschitz condition is a smoothness condition for functions and it is stronger than regular continuity. Explicitly, Lipschitz condition says for function $f(t, x)$, if (t, x) and (t, X) are two points in the domain D , then

$$|f(t, x) - f(t, X)| < K|x - X| \quad (\text{A.5})$$

where K is the usually called Lipschitz constant of the function f and $|\cdot|$ represents some distance measurement. Ince showed that the continuity of $f(x, t)$ is not necessarily and the requirements for $f(x, t)$ are that it is bounded and all the integral of

the form $\int_0^t f(x^i, s)ds$ exist [73].

For a system equation having the following form, with m dependent variables and each of the equation has a first order form

$$\frac{dx_1}{dt} = f_1(t, x_1, x_2, \dots, x_m) \quad (\text{A.6})$$

$$\frac{dx_2}{dt} = f_2(t, x_1, x_2, \dots, x_m) \quad (\text{A.7})$$

$$\vdots \quad (\text{A.8})$$

$$\frac{dx_m}{dt} = f_m(t, x_1, x_2, \dots, x_m) \quad (\text{A.9})$$

Ince [73] showed that if f_1, f_2, \dots, f_m are single-valued and continuous in the $m + 1$ variables which are restricted to lie in the domain D defined by

$$|t - t_0| \leq a, |x - x_0| \leq b_1, \dots, |x - x_m| \leq b_m \quad (\text{A.10})$$

if the greatest of the upper bounds of f_1, f_2, \dots, f_m in the domain D is M ; if h is the least of $b_1/M, b_2/M, \dots$, and b_m/M ; let t be further restricted if necessary by $|t - t_0| < h$; the Lipschitz condition has a form as

$$\begin{aligned} & |f_r(t, X_1, X_2, \dots, X_m) - f_r(t, x_1, x_2, \dots, x_m)| < \\ & K_1|X_1 - x_1| + K_2|X_2 - x_2| + \dots + K_m|X_m - x_m| \end{aligned} \quad (\text{A.11})$$

where $r = 1, 2, \dots, m$; the iteration having the form as

$$x_r^i(t) = x_r(t_0) + \int_0^t f_r(x_1^{i-1}, x_2^{i-1}, \dots, x_m^{i-1}, s)ds \quad (\text{A.12})$$

will converge to the unique and continuous solution of Eq. A.6 to A.9 .

APPENDIX B

CHEBYSHEV POLYNOMIALS

Chebyshev polynomials are a sequence of orthogonal polynomials developed by the Russian mathematician Pafnuty Lvovich Chebyshev in 1857 [74]. There are two kinds of Chebyshev polynomials. The k^{th} Chebyshev polynomials of the first kind usually are denoted by T_k and the k^{th} Chebyshev polynomials of the second kind usually are denoted by U_k . Through the dissertation, wherever there exist no confusions, we call Chebyshev polynomials of the first kind as Chebyshev polynomials for simplicity.

The Chebyshev polynomials can be defined through the recurrence relation as

$$T_0(\tau) = 1 \quad (\text{B.1})$$

$$T_1(\tau) = \tau \quad (\text{B.2})$$

$$T_{k+1}(\tau) = 2\tau T_k(\tau) - T_{k-1}(\tau) \quad (\text{B.3})$$

or the Chebyshev polynomial of degree k can be defined by using trigonometric identity

$$T_k(\tau) = \cos(k \arccos(\tau)), \tau \in [-1, 1] \quad (\text{B.4})$$

The continuous orthogonality of Chebyshev polynomials satisfies

$$\int_{-1}^1 T_n(x) T_m(x) \frac{dx}{\sqrt{1-x^2}} = \begin{cases} 0 : n \neq m \\ \pi : n = m = 0 \\ \frac{\pi}{2} : n = m \neq 0 \end{cases} \quad (\text{B.5})$$

The discrete orthogonality of the Chebyshev polynomials using the CGL nodes

satisfies

$$\sum_{k=0}^{k=N} {}''T_n(x_k)T_m(x_k) = \begin{cases} 0 : n \neq m \\ N : n = m = 0 \\ \frac{N}{2} : n = m \neq 0 \end{cases} \quad (\text{B.6})$$

where $''$ represents that both the first and last terms in the summation are halved, and the $(N + 1)$ CGL nodes for the N^{th} order Chebyshev polynomials are calculated from

$$x_k = \cos\left(\frac{k\pi}{N}\right) \quad (\text{B.7})$$

The integration of the Chebyshev polynomials has the property

$$\int T_k(x)dx = \frac{1}{2}\left(\frac{T_{k+1}}{k+1} - \frac{T_{k-1}}{k-1}\right) \quad (\text{B.8})$$

and the first derivative of the Chebyshev polynomials of the first kind is related with the Chebyshev polynomials of the second kind through

$$\frac{dT_k}{d\tau} = kU_{k-1} \quad (\text{B.9})$$

Fox [18] proved that if a continuous function $f(\tau)$ is approximated by a n^{th} order Chebyshev polynomials as

$$f(\tau) \approx p_n(\tau) = \sum_{k=0}^n \alpha_k T_k(\tau) \quad (\text{B.10})$$

and if the coefficient α_k is calculated by

$$\alpha_k = \frac{2}{N} \sum_{j=0}^N {}''f(\tau_j)T_k(\tau_j), k = 0, 1, \dots, n \quad (\text{B.11})$$

where the $(N + 1)$ discrete nodes are chosen by

$$\tau_j = \cos(j\pi/N), j = 0, 1, \dots, N \quad (\text{B.12})$$

the error $e_n(\tau) = f(\tau) - p_n(\tau)$ satisfies the discrete least-squares criterion

$$S = \sum_{j=0}^N e_n^2(\tau_j) = \textit{minimum} \quad (\text{B.13})$$

and

$$S_{min} = \sum_{j=0}^N (f^2(\tau_j) - \sum_{r=0}^N \alpha_r^2 T_r^2(\tau_j)) \quad (\text{B.14})$$

In the previous equations, Σ' denotes that the first term is halved and Σ'' represents that the first and last terms are halved.

Compared with the interpolation formulae presented by Fox [18], these discrete formulas from Eqs. (B.10) to (B.12) only require the order of polynomials n be equal or less than the number of nodes N , which offers some flexibility for choosing nodes and polynomial orders. The explicit expression also provides an error bound in Eq. (B.14) which is difficult to obtain by using the interpolation formulae. Also notice there exists another set of Chebyshev nodes defined by

$$\tau_j = \cos\left(\frac{2j+1}{N+1} \frac{\pi}{2}\right), j = 0, 1, \dots, N \quad (\text{B.15})$$

which are essential the zeros of Chebyshev polynomials, also provide a formula for discrete least-squares fit as

$$p_n(\tau) = \sum_{k=0}^n {}' \beta_k T_k(\tau) \quad (\text{B.16})$$

with

$$\beta_k = \frac{2}{N+1} \sum_{j=0}^N {}'' f(\tau_j) T_k(\tau_j) \quad (\text{B.17})$$

Fox compared these two different formula and states that the first set has theoretical advantages for slow convergence cases, and also is economic as the old matching points can be reused when the number of points doubles [18]. We choose to use Eqs. (B.10) to (B.12) in this dissertation and require $n \equiv N$ for simplicity.

The first six orders of the Chebyshev polynomials are shown in Fig. 94.

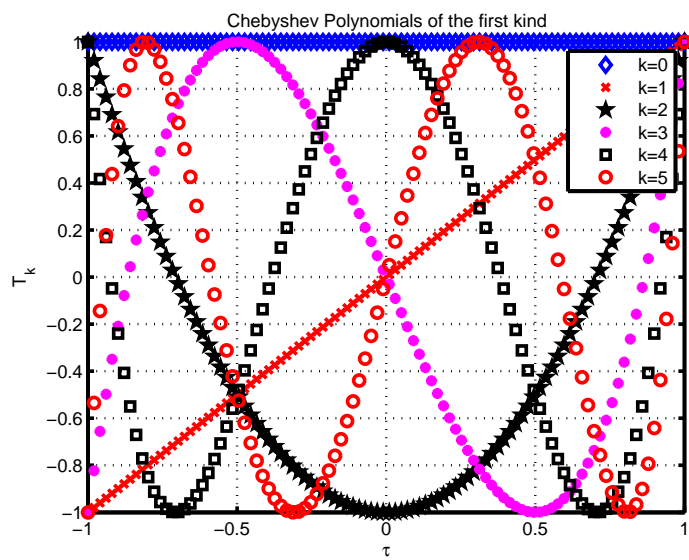


Fig. 94. Chebyshev Polynomials of the first kind

VITA

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