A GRAPH THEORETIC APPROACH AND SOUND ENGINEERING PRINCIPLES FOR DESIGN OF DISTRICT METERED AREAS Giada Ferrari¹, Dragan Savic², Gianfranco Becciu³

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5 ABSTRACT

6 The design of District Metered Areas (*DMAs*) in existing water distribution networks, 7 especially in urban areas, involves a high number of decision variables and the effects of 8 implementing them in districts have to be evaluated, in order not to affect the quality of the 9 service to customers.

10 A new methodology for designing a given number of districts in looped water 11 distribution networks, is proposed here. It is based on graph theory and takes into account 12 some important *DMA* design criteria: the maximum and minimum size recommended for a 13 district, the connectedness of each district to the water supply source and the absence of links 14 between the districts: therefore it allows the creation of *DMA*s that are independent one from 15 another.

16 A recursive bisection procedure has been applied to create districts, while an 17 algorithm for graph traversal has been used to verify whether each district can be reached 18 from the water source and connectivity between the nodes.

- The successful application of the proposed methodology to a case study has proved its
 effectiveness for District Metered Areas design in real urban water distribution networks.
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22 INTRODUCTION

Water distribution networks, especially in urban areas, are usually designed as looped systems, where water can flow in different directions and demand nodes can be served through different paths. In such systems multiple flow paths prevent stagnation from occurring, allow water pressure to be uniform, and provide a high level of reliability to the supply service, because the demand nodes can be served even in the case of pipe failures. For these reasons, looped distribution networks are preferred for urban water supply.

29 However, sometimes it may be convenient to divide the network into district metered 30 areas (DMAs), that are areas in the water network, independent from each other, created by 31 the closure of valves or disconnection of pipes and where the inlet and outlet flow is metered 32 (Figure 1). As a matter of fact, experience has shown that the implementation of district 33 metered areas provides a series of benefits for water distribution systems: for example it 34 allows for the reduction of leakages, due to easier and faster identification and location of 35 leaks, and for the creation of a permanent pressure control system, which enables a low level 36 of leakage to be maintained (AWWA 2003). It also improves the water distribution network 37 management thanks to the simplified evaluation of the water balance and reduces water 38 security risks, since the potential movement of contaminants throughout the system is 39 minimized.

Since *DMAs* have been introduced in water distribution systems, considerable work has been done to improve the design, planning and management of *DMAs*: indeed, a number of studies and technical reports containing guidelines, design criteria and recommendations are available in literature (see e.g. Farley 1985, Morrison et al. 2007, Baker 2009). For example, indications are given about the minimum and maximum number of customer connections that a district should contain (*DMA*'s size); the main transmission system should be kept separated from the *DMAs* in order to ensure a flexible and reliable water supply; each

DMA should be connected directly with the transmission main and be independent, i.e.
without any connection with other *DMAs*; other factors to take into account are pressure
constraints at demand nodes, final leakage level target, implementation and maintenance
costs.

51 The re-design of a water distribution network into DMAs is not a trivial issue and if 52 not undertaken with care, can lead to supply problems, reduction of reliability and worsening 53 water quality (Grayman et al. 2009). Due to its complexity, the design of DMAs in practice 54 has always been made empirically: an initial division into districts, each having appropriate 55 size, is considered; then hydraulic simulations are performed under different demand 56 scenarios (average demand, maximum demand, fire flow, pipe burst) and modifications are 57 performed manually in the water network if pressure constraints are not verified. In other 58 words, an iterative trial and error approach is adopted until an acceptable solution is found.

Murray et al. (2010) compared the performances of looped water distribution networks with the corresponding districted version. They showed how, if the criteria mentioned above are followed, the creation of *DMA*s brings advantages in terms of water security and leakage reduction, without compromising either the reliability or the quality of the water supplied. Performance diminishes when connections between *DMA*s cannot be avoided and water from some districts flows into downstream *DMA*s, emphasizing the importance of the independence between them.

Recently, some new methods for designing *DMAs* in existing water distribution networks have been developed. Award et al. (2009) used genetic algorithms for individuating the optimal setting of Pressure Reducing Valves (*PRVs*) in water distribution networks. The methodology also involved finding the optimal *DMA* boundaries such that the excessive outlet hydraulic pressure at nodes can be minimised throughout the day. They adopted a fitness function that assesses the annual benefits deriving from the implemented pressure

management scheme, and therefore the cost saving. The water network is considered as a graph and the creation of individual solutions in initial population is made by applying the Depth First Search algorithm (Cormen et al. 2001). A spanning tree is grown from the water source, then a certain number of PRVs are placed on tree branches at random and the remaining links are reinstated. This procedure ensures the connectivity and reachability of each demand node.

Di Nardo and Di Natale (2009) developed a decision support system for *DMA* design using a graph theory approach based on a shortest path algorithm (Dijkstra 1959). Their method is based on the determination of a set of candidate pipes to be closed in order to form the boundaries of the districts. The minimum head loss paths from the water source to each demand node in the network are found and each pipe is associated with a frequency value, which is proportional to the number of times that the pipe is found in a path. Candidate pipes are those having a frequency lower than a certain threshold fixed by the decision maker.

85 The same authors (Di Nardo and Di Natale 2011) applied the graph partitioning 86 methodology by Karapys and Kumar (1995) to the DMAs design problem. They achieved the 87 division of network nodes into a certain number of districts (partition), approximately equally 88 sized, so that the number of links to be closed, also called cutting edges, was minimized. The 89 methodology required that weights related to hydraulic properties were assigned to links and 90 vertices of the graph. Pipe flow, dissipated power or diameters were chosen as hydraulic 91 properties assigned to the links while water demand was associated to the vertices. The 92 choice was made by analyzing all the possible combinations among the boundary pipes found 93 with the graph partitioning algorithm.

Alvisi and Franchini (2012) proposed an automatic procedure based on the generation of different *DMAs* network layouts that are further compared in order to find the most resilient one. The creation of a solution, i.e. a possible *DMA* layout, is made with the aid of

97 the Breadth First Search algorithm (Pohl 1989): from each node of the network a tree is made 98 to grow and all nodes are grouped in sets, according to their distance from the source node, 99 and the cumulative water demand is evaluated. Districts are defined as groups of sets having 100 total water demand within the recommended limits.

Di Nardo et al. (2013) recently developed a methodology for water network sectorization that allows to divide the water network into isolated areas (sectors) that are completely separated from one another and fed by their own water source. The authors employed graph theory to model the water system and applied a heuristic optimisation technique aimed to minimise the dissipated power in order to determine the sectors' boundaries.

107 Lastly, Diao et al. (2013) determined DMAs boundaries through an approach based 108 on the assumption that water distribution systems can be seen as 'community structures'. A 109 'community' is defined as a group of nodes characterised by a high density of edges between 110 its vertices, whereas the number of edges connecting a community with another is 111 significantly lower. The authors identified communities in a water distribution network 112 starting from a condition in which each single vertex is a community and adopting a greedy 113 strategy, joining communities until an increase in the modularity, an indicator of how well a 114 graph is divided into communities, is observed. The size of the communities was then 115 calculated and those exceeding the upper limit for DMA size were further decomposed, in 116 order to ensure that each community had a suitable number of customer connections. Finally, 117 feed lines, i.e. pipes connecting communities, and pipes to be closed were identified such that 118 the minimum pressure requirements at each node are verified.

The aforementioned studies are based on graph theory principles and techniques. This means that the water distribution network is handled as a graph, whose vertices are represented by demand nodes, reservoirs and tanks, while edges are represented by pipes,

122 pumps, and any other link between two vertices in the water system (Deuerlein 2008, 123 Kesavan 1972). As a matter of fact, graphs have been proved to represent a useful and 124 powerful tool for modelling a water distribution network, especially when dealing with the 125 DMA's design problem (Ostfeld and Shamir 1996, Ostfeld 2005, Tzatchkov 2006). 126 Furthermore, graphs are quite simple structures, so the complexity of the problems can be 127 reduced significantly. Of particular interest, regarding the design of districts in a water 128 distribution network, are the properties of connectivity among nodes within a DMA and 129 reachability of each node from the water supply source. The former indicates that all the 130 nodes belonging to the same district are connected, while the latter that all the nodes in the 131 network have a direct path either to the water source, or to the main transmission system.

132 However, water distribution networks are not just graphs composed by arcs and 133 vertices. There are functional requirements that need to be satisfied, such as the delivery of a 134 certain amount of water to each node and the maintenance of adequate pressures. Moreover, 135 flow directions depend on hydraulic laws, in particular on the continuity equation at nodes 136 and the energy conservation around each elementary loop. Therefore, when using graph 137 theory for water distribution network analysis, the hydraulic and connectivity constraints 138 have to be taken into account. The former ensures the physics of the water flow is taken into 139 account, and the latter makes sure that no nodes are isolated, and thus all the customers in the 140 water supply distribution networks are properly served.

As mentioned above, when designing *DMA*s in an existing water distribution network some criteria have to be followed, in order to not lower the performance of the water supply system. The essential factors to take into account are the size limits of the *DMA*s, the connectivity properties and respecting the minimum pressure constrains. Nevertheless, the approaches developed in literature, although able to determine *DMA*'s boundaries, do not take into account all the aforementioned factors in the design process. Either the methodology

147 by Award et al. (2009) or the one by Di Nardo and Di Natale (2009) consider the size of the 148 resulting DMAs as a design criterion: the first aims only to minimise the excess of pressure at 149 nodes, and the second chooses the pipes to be closed only on the basis of minimum head loss 150 paths from the water source to each demand node. Moreover, the majority of the approaches 151 illustrated cannot ensure the resulting DMAs be supplied independently. Non-independent 152 DMAs and DMAs which lack a direct connection to the transmission mains are likely to 153 occur, despite the shortcomings in terms of reliability and water quality this would produce. 154 Only the methodology developed by Di Nardo et al. (2013) allows to identify DMAs that are 155 independent, but the authors did not consider the recommendations about the number of 156 customer connections per district as a design criterion.

157 This paper presents a methodology based on graph theory that can partition an 158 existing water distribution network in DMAs. Its purpose is to determine the boundaries of 159 DMAs ensuring that each DMA be between the recommended size limits, and be directly 160 connected to the main transmission system. The methodology will also ensure that no links 161 exist between different DMAs, thus there is no flow exchange between adjacent DMAs, and 162 all the demand nodes are adequately served during the simulation period. This paper will also 163 show how it is able to divide a water network into a certain number of DMAs which are 164 spaced at specified intervals and with the important characteristic of being independent of 165 one another.

166

167 THE METHODOLOGY

168 The methodology developed, represented in Figure 2, consists of three main steps 169 (corresponding to the square boxes in the flow diagram):

- A preliminary analysis of the water distribution network, to identify the transmission
 mains, the "independent" districts (see Section 1 below) and the number of *DMAs* to
 create. *k*.
- A recursive bisection algorithm, which has been tailored to take into account the design criteria described in the previous section. It determines the *DMA*'s boundaries, where valves are assumed to be available, and therefore the pipes that have to be closed. If, in practice, valves were not available at the boundaries, the expenditure for their installation would need to be taken into account in the total cost of *DMA*s implementation.
- A hydraulic simulation to verify that minimum pressure requirements are satisfied at
 each node.
- 181 This section aims to provide a detailed analysis of the steps listed above, along with a182 description of the output.
- 183 1. Input data
- 184 The proposed methodology requires the following input data:
- the model of the water distribution network, in particular the topology, i.e. the spatial
 coordinates of nodes and links between them, the hydraulic characteristics of the
 network, among which nodal base demands, demand patterns, head loss formula,
 pipes roughness, etc., and the hydraulic analysis method, for instance the gradient
 method (Todini and Pilati 1987)
- minimum required pressure to ensure the delivery of water to the customers, to verify
 the hydraulic feasibility of the resulting network
- the minimum and maximum allowable size for a single DMA, C_{min} and C_{MAX} , usually expressed in terms of the number of customer connections per district and typically equal to 500 and 5000 respectively (Farley 1985, Morrison et al. 2007). In those cases

when the information about the number of connections at each node is not available,
the minimum and maximum size of a *DMA* can be expressed in terms of total water
demand within the *DMA*.

198 An average relationship between the number of connections and water demand is 199 adopted: let C_{tot} be the total number of customer connections in the network and W_{tot} the total 200 average water demand (l/s); $W_{d,min}$ and $W_{d,MAX}$, respectively the minimum and maximum 201 average water demand of a generic independent district, represent the minimum and 202 maximum size per district and are defined as follows:

203
$$W_{d,min} = \frac{W_{tot}}{c_{tot}} C_{MAX}$$
 Eq. 1a

204
$$W_{d,MAX} = \frac{W_{tot}}{c_{tot}} C_{min}$$
 Eq. 1b

The proposed methodology allows only for the creation of *DMA*s whose sum of nodal water demands is in between the boundary values given in Equations 1a and 1b.

207 2. Preliminary analysis

208 A preliminary analysis of the water distribution network to be divided into *DMA*s is 209 necessary, in order to acquire some basic information, fundamental for performing the further 210 steps of the methodology. At first, water transmission mains have to be identified: they are 211 large water pipes used to extend and convey water between sources, such as storage facilities, 212 facilities, external water supply networks, wells, springs, etc., and the distribution mains, and 213 are not intended to serve as a distribution system. The distinction made in the definition of 214 transmission mains and distribution mains is a function of the size of the water system itself. 215 Thus, this distinction is not universal and varies from system to system, and the actual 216 function of the mains in the system is more important than the size in determining what 217 category a main falls under. According to usual practice, transmission mains can be

considered as the series of connected pipes having a diameter equal to or greater than a
threshold of approximately 300 - 350 mm (12 - 14 inches).

220 When the transmission mains have been identified, the network is analysed in order to 221 find the "independent districts", i.e., groups of nodes linked to the transmission main and 222 having no connections with any other group (an example of independent districts is 223 represented in Figure 3). This analysis is performed by exploring the water distribution 224 network with the Breadth First Search (BFS) algorithm. All the independent districts are 225 automatically detected by considering every node that belongs to the transmission mains as 226 sources from which to start the exploration of adjacent nodes. From each of them, the 227 algorithm, in its first stages, detects all the nodes that are at the distance of one edge from the 228 source and store them in a list. In the second stage, all the nodes placed at a distance of two 229 edge from the source are found and added to the list. The BFS algorithm terminates when 230 there are no more nodes reachable from the source. These steps are performed for each node 231 of the transmission mains system. The order in which nodes are explored does not influence 232 the result of the algorithm, i.e. the list of nodes reachable form each source. Some of these 233 lists might contain the same nodes, because it can happen that a group of nodes can be 234 reached from more than one source. This group, having a number of links with the 235 transmission mains equal to the number of sources from which its nodes can be explored, 236 represents an independent district.

With the aid of BFS algorithm, all the independent groups of nodes are easily and
quickly identified, along with their connections with the main transmission system; their size
in terms of total nodal water demand is also calculated.

The dimensions and the characteristics of these groups of nodes can vary considerablyand depend on the water distribution systems being analysed.

- Specifically, the size, i.e. the number of customer connections C_d of each independent
 district, can be:
- lower than the minimum acceptable for a single DMA (C_d < C_{min}),
- between the minimum and the maximum ($C_{min} < C_d < C_{MAX}$),
- 246
- greater than the maximum ($C_d > C_{MAX}$).
- A similar distinction can be made with respect to the water usage, as seen in equation
- 248 1.

Such a characterization allows for the location of *DMA*s that might already exist in the water distribution network, that are groups of connected nodes having size between the limits, that are 500 and 5000 customer connections in this study.

252 In the first case ($C_d < C_{min}$), the independent district is too small to be considered a 253 DMA. Therefore, these nodes keep being supplied directly from the transmission mains, no 254 meters are installed and no modifications are made to the network. Conversely, if the size is 255 between the boundary values C_{MAX} and C_{min} given as input, the independent district is 256 considered as a DMA. Therefore during the preliminary analysis some DMAs might be 257 identified. Generally, it is not necessary to insert isolating valves, because there are no 258 connections with any other district (i.e., the DMA is already isolated from the rest of the 259 network), and a meter is installed on the feeding pipe that link the district with the 260 transmission main. However, if the district is fed from multiple sources, further analyses need 261 to be carried out to determine whether it is convenient to set valves on one or more feeding 262 pipes. Lastly, if the size of the independent district is greater than the maximum allowable for 263 a single DMA, it needs to be divided into smaller districts.

The preliminary analysis ends with determining the number k of *DMA*s to create in each independent district belonging to the third group ($C_d > C_{MAX}$). This number is chosen between a minimum and maximum value, depending on the minimum and maximum size

267 allowable. Nevertheless, the number of DMAs that can be created within an independent 268 district, cannot exceed the number of pipes linking the district itself with the transmission 269 main, because the methodology aims to create DMAs all having a direct path to the water 270 source, as recommended by the design criteria. In the case there were only few connections 271 between the independent district and the transmission mains, an appropriate number of 272 linking pipes should be added to the existing network, in order to allow the creation of a 273 certain number of DMAs. However, this situation is very unlikely to occur, because the 274 supply of an independent district having such a high number of customer connections could 275 hardly come from a limited number of feeding pipes from the transmission main system.

Therefore, if n_c is the number of connections with the transmission main, the minimum and maximum number of district are given by equations 2 and 3:

278
$$k_{min} = \frac{W_d}{W_{d,MAX}}$$
 Eq. 2

279
$$k_{MAX} = \min\left(\frac{W_d}{W_{d,MAX}}, n_c\right)$$
 Eq. 3

Since small *DMAs* have high implementation and maintenance costs whilst large *DMAs* do not provide as much benefit in terms of water network control and management, *k* should be chosen close to the mean value.

283 3.

The recursive procedure

Given a certain value of k, the *DMA*'s boundaries are defined through the application of a recursive bisection algorithm. It is based on the principle of recursively bisecting the set of demand nodes until the required number of subsets is obtained. The procedure for performing a single bisection, illustrated in the following section, is crucial to achieve the desired result, i.e., a feasible water distribution network that is divided into a certain number of subsets (*DMAs*) having characteristics that reflect the design criteria. The diagram in Figure 4 illustrates the recursive procedure.

Let SET={ $n_1, n_2, n_3, ..., n_n$ } be the set of nodes in a graph that are required to be divided into a certain number of subsets n_{SS} , equal to or greater than 2. When the graph is a water distribution network, SET corresponds to the set of demand nodes, the subsets are the *DMAs* and n_{SS} is equal to *k*, set by the decision maker between the minimum and the maximum value previously evaluated. The first bisection produces two subsets A and B. The number of subsets to be further created from A and B is given by rounding down half of n_{SS} and rounding up half of n_{SS} respectively (equations 4 below).

298
$$n_{SS,A} = floor\left(\frac{n_{SS}}{2}\right)$$

$$n_{SS,B} = ceil\left(\frac{n_{SS}}{2}\right)$$
Eq. 4

299 Since n_{ss,A} and n_{ss,B} can have differing values, the size of A and B can be different as 300 well: for instance, if the number of subsets to be created from A is higher, the size of A has to 301 be greater than the size of B. The size of a subset, for water distribution networks, is 302 represented either by the total number of customer connections within the subset or by the 303 water use within the subset. Since the number of customer connections per district is one of 304 the factors to be taken into account in the DMA's design, the size of a subset has to be defined 305 appropriately. A correct definition of the subset's size leads to DMAs whose number of 306 customer connections satisfies the recommended design criteria.

Let S_A and S_B be the sizes of the first two subsets, respectively A and B; their values
have to be between the boundaries given in equations 5.

309
$$S_{A/B,min} = n_{ss,A/B} W_{d,min}$$
$$S_{A/B,MAX} = n_{ss,A/B} W_{d,MAX}$$
Eq. 5

The definition of the subset's size as in these equations has two main purposes: dividing the network into a number of districts, k, having roughly the same size and following the guidelines given in literature. However, it is not recommended that the subset size is too close to the boundary values, especially if $n_{ss,A}$ and $n_{ss,B}$ are greater than two. In fact, a size

314 value very close to the boundary could lead to final *DMAs* having vastly different 315 dimensions, implying a higher number of valves to be closed and thus a higher cost. An 316 acceptability range for the size of subsets A and B is then defined (Figure 5) and its boundary 317 values are given by equations 6 - 8:

318
$$S_{A/B,low} = al \ n_{ss,A/B} (W_{A/B,med} - W_{d,min})$$
 Eq. 6

319
$$S_{A/B,up} = al \ n_{ss,A/B} (W_{d,MAX} - W_{A/B,med})$$
Eq. 7

320
$$W_{A/B,med} = n_{ss,A/B} \frac{\sum_{i=1}^{n} W_i}{k}$$
 Eq. 8

321 where *al* is a parameter between 0 and 1 that indicates the width of the interval around 322 the medium water usage value and allows for flexibility in the DMA's size. The choice of al 323 is arbitrary and different values can be set for obtaining different *DMA* layouts and to check 324 the sensitivity of the result with respect to the parameter. If *al* is set equal to 1, the upper and 325 lower bounds for the subset size are respectively the maximum and minimum allowable. As 326 previously mentioned, the adoption of a restricted admissible interval for subset size proves 327 to be particularly useful when k is high (greater than 5), since this provides for approximately 328 equally sized districts.

329 After the first bisection has been performed and subsets A and B have been 330 determined, if $n_{ss,A}$ is equal to or greater than 2, A becomes the new set to be bisected and the 331 previous steps are repeated until $n_{ss,A}$ is lower than 2. The same is done with the subset B.

332 The steps of the recursive bisection, illustrated in Figure 4, are summarized below:

333 1. Given SET=
$$\{n_1, n_2, n_3, ..., n_n\}$$
 and $n_{ss} = k_{ss}$

334 2. Evaluate $n_{ss,A}$ and $n_{ss,B}$ with equations 4;

335 3. Bisect SET into two subsets A and B such that their size S_A and S_B are between S_{A,low}
336 and S_{A,up} and S_{B,low} and S_{B,up} respectively

337 4. a) If $n_{ss,A} \ge 2$ update SET = A, $n_{ss} = n_{ss,A}$ and go back to step 2)

338

b) If $n_{ss,B} \ge 2$, update SET =B, $n_{ss} = n_{ss,B}$ and go back to step 2)

- 339 5. Stop when both $n_{ss,A}$ and $n_{ss,B}$ are lower than 2.
- 340 4. The bisection algorithm

341 The bisection algorithm presented here is the key part of the whole methodology. By 342 its recursive application, it allows for the creation of districts having appropriate size, i.e. the 343 total water demand is between the limits given in equation 1, and that are independent of each 344 other. This independence ensures that each DMA is connected to the transmission mains by at 345 least one pipe, so that at least one direct flow path exists between the water source and each 346 demand node, and there are no flow exchanges between contiguous DMAs. A random 347 component has been included in the process defining DMAs, so that each single run of the 348 algorithm produces a different layout of the DMA where all the districts have at least one 349 connection with the transmission main. This capacity to provide a number of alternative 350 solutions is particularly useful because the decision maker is provided with a number of 351 feasible options. In fact, when dealing with large water distribution systems, there are many 352 possible divisions of the network into districts. The decision maker can then choose the 353 division that best suits his requirements from the various options provided.

Figure 6 shows the bisection algorithm that is applied recursively to identify the desired number k of districts in the water distribution network. Figure 7 shows an example of an independent district in a real large water distribution network. Nodes belonging to SET are highlighted in grey, nodes belonging to C in black, and the transmission mains are indicated with a black bold line. The adjacency matrix, which is an n-by-n matrix that embodies the information about the connections between the nodes belonging to SET, is determined.

360 Let SET= { id_1 , id_2 , id_3 , ..., id_n } be the set containing the n nodes belonging to an 361 independent district having size greater than C_{MAX}, i.e. the set of nodes that need to be

- 362 grouped into districts; let $C = \{c_1, c_2, c_3, ..., c_{nc}\}$ be the set of the n_c nodes of SET adjacent to
- the transmission main.
- 364 The procedure for performing a bisection of a certain set of nodes SET, illustrated in
- 365 Figure 6 in flow chart form, consists in the following steps:
- 366 1. A node $c_i \in C$ is chosen at random
- 367 2. A spanning tree is made to grow from c_i using the BFS algorithm
- 368 3. Let $L = \{c_i, n_1, ...\}$ be the list of nodes in the order they are discovered with the BFS:
- 369 since all the nodes in SET are connected, they will all be discovered sooner or later, L
- is thus an n-dimensional array
- 371 4. Evaluate the cumulative water usage for the set L, $W = \{W_1, W_2, ..., W_n\}$
- 372 5. Let I be the set of nodes n_j such that $S_{A,low} < W(n_j) < S_{A,up}$ (see Figure 8)
- 373 6. Choose a node $n^{I} \in I$ at random
- 374 7. Let SET1 be a subset of L formed by all the nodes in the list between the tree source 375 c_i to n^I : SET1 = { $c_i, n_j, ..., n^I$ }
- 376 8. Let SET2 be the complement of L with respect to SET1: SET2 = SET SET1
- 377 9. Consider the set C'= C \subset SET2, containing the nodes of SET2 adjacent to 378 transmission main
- 379 10. Grow a tree T_k from each $c_k \in C'$ and evaluate their corresponding water usages W_k
- 380 11. a) If there is a W_k such that $S_{B,min} < W_k < S_{B,MAX}$, set $B = T_{k}$, $A = SET T_{k}$, 381 otherwise
- 382 b1) update set I: $I=I \{n^I\}$; if I is not empty, go back to 6), otherwise
- 383 b2) $C = C \{c_i\}$ and go back to 1)
- 384 12. Evaluate the number of connections that A and B have with the transmission mains:

385 $n_{C,A}=C \cap A, n_{C,B}=C \cap B$

386 13. If $n_{C,A} \ge n_{ss,A}$ and $n_{C,B} \ge n_{ss,B}$, A and B are the result of the bisection, otherwise go 387 back to 11.b1).

The last step verifies that the number of pipes connecting A and B with the transmission main is equal to or greater than the number of subsets to be further made from A and B, because every *DMA* must have a direct path to the water source. Therefore each district that is to be created will have a direct connection with the transmission main and the independence of every *DMA* is ensured.

393 5. Outputs

The procedure illustrated above allows for the division of an independent district into a certain number of smaller districts that meet the size and connectedness requirements. It gives as its output the list of nodes belonging to each subset, that is the list of nodes belonging to each *DMA*.

398 Its application to all the independent districts provides a possible *DMA* layout for the 399 water distribution network under examination. The integration of a random component in the 400 bisection algorithm (steps 1 and 6) allows for each run of the procedure to generate a 401 different solution.

402 The solution obtained describes where a certain number of pipes have been closed in 403 order to create the DMA's boundaries in the original/starting water distribution network. 404 Hence, the hydraulic feasibility of the solution found needs to be checked. A hydraulic 405 simulation is then performed using the software EPANET and the pressures at each demand 406 node are evaluated. The last step is to verify that each demand node during the whole 407 simulation period is associated a hydraulic head equal to or greater than the minimum 408 required. If this constraint is not verified, then the solution is discarded as unfeasible, and a 409 new division onto the DMAs is sought until one that respects the minimum pressure 410 requirements is determined. Here, the ability of the resulting system to meet fire flow

411 conditions has not been checked. However, the creation of DMAs has been proven to affect 412 the capability of meeting fire flow conditions, as well as the reliability and the quality of the 413 water delivered. Therefore, further analysis will be required to determine the best solution 414 among those generated in relation to these three aspects.

The value of k chosen by the designer can also be changed in order to find a wider variety of possible solutions. A number of solutions can be found and compared on the basis of performance indicators, in order to determine which one is the best to adopt in a specific water network.

419

420 CASE STUDY

The effectiveness of the proposed methodology was proven by applying it to a real case study. The water distribution network used is a modified real world system, frequently used as a test bed for various modelling exercises including the Battle of the Water Sensor Networks competition (Ostfeld et al. 2008). Its geographic representation and the names of the components have been distorted in order to protect the identity of the system. The adjustments made are related only to the appearance and do not have any influence on the connectivity and the hydraulic behaviour of the network.

The network serves approximately 150,000 people and includes two reservoirs, two tanks, four pumps and five valves, and it represents a typical example of a looped urban distribution system. Its topology is illustrated in Figure 9.

The hydraulic analysis was performed by using the simulation software EPANET (Rossman 2000). The head losses were evaluated with the Hazen-Williams formula, while other basic characteristics of the network that has been used in the hydraulic analysis, are reported in Table 1.

435 The minimum and maximum number of customer connections per district were set 436 equal to 500 and 5000 respectively. Nevertheless, the data about the number of connections 437 per node was unknown, and the only information available was the total number of customer 438 connections in the system, that is equal to 77916. Therefore, an average relationship between connections and water demand was used (Gravman et al. 2009): let \overline{W}_c be the average water 439 440 use per connection, given by the ratio between the total average water use and the number of 441 connections in the network. The upper and lower limits for the size of a district can be 442 expressed in terms of water usage, W_{d,MAX} and W_{d,min}, and can be evaluated as the product 443 between the average water use per connection and the upper and lower limit of connections 444 (equations 1a and 1b). The resulting values are reported in Table 2, along with the variables 445 used for their evaluation.

The first step of the preliminary analysis of the water distribution network, whose purpose is to gain detailed information about the water system, is to identify the pipes belonging to the transmission mains. In this case study, the main transmission system is considered to be composed of all the connected pipes having a diameter greater than or equal to 350 mm (14 inches). Transmission mains are shown in Figure 9. They are composed of 870 pipes, and represent 9.2% of all the pipes in the network, totalling 162.86 km in length.

452 The second step is to identify the independent districts in the water distribution 453 network, that are groups of fully connected nodes having at least one link between the district 454 and the transmission main system (that is a direct connection with the water source). The 455 independent districts are then classified according to their size, i.e. according to the total 456 water use, that is the sum of the water demands of nodes within the district. This step is 457 performed by exploring the network as a graph with the Breadth First Search (BFS) 458 algorithm: all the nodes belonging to the transmission main are eliminated from the graph, 459 i.e. the values of the corresponding rows and columns in the adjacency matrix are set equal to

460 zero. Their neighbour nodes are instead considered and set as sources from which to start 461 growing a spanning tree performing BFS (but also Depth First Search can be employed). The 462 list of nodes explored from each "source" represents an independent district, i.e. a group of 463 connected nodes linked by at least one pipe to the transmission main.

The total water use W_d of each independent district is then evaluated: groups of nodes characterized by $W_d < W_{d,min}$, that have water use lower than 8 l/s, are too small to be considered *DMA*s, and will be ignored in the further analysis. The total number of this type of nodes is 946, and the corresponding total water demand is 125.3 l/s; they thus represent 9.3% of the total water demand of the system and their location on the map is shown in Figure 10.

469 Groups of nodes characterized by W_d between W_{d,min} and W_{d,MAX} are considered as 470 DMAs, because their size respects the requirements given in Table 2: they are districts 471 already existing in the network, that do not require any valve to be inserted to define their 472 boundary: only a meter on the pipe linking them with the transmission main will be installed 473 in order to measure the inlet and outlet flows. In the water system in analysis there are 20 474 districts of this type, they are illustrated in Figure 11 and their characteristics are summarized 475 in Table 3. The total water demand of the nodes belonging to the these DMAs is equal to 476 514.9 l/s, which corresponds to the 34.06% of the total water demand of the system.

Finally, groups of nodes having total $W_d > W_{d,MAX}$ are identified (Figure 12): there are three of them in the water network, and their total water use is 819.9 l/s that is the 54.23% of the water demand of the entire network. Each of these three "large districts", whose characteristics are summarized in Table 4, will be divided into smaller *DMA*s applying the developed methodology; the solutions obtained will subsequently be combined in order to define a number of alternative solutions for the whole network, for which the performances will be evaluated.

The minimum and maximum number of districts that can theoretically be created in each one of the three large independent districts, depending on their water use and on the size limits reported in Table 2 are evaluated by the equations 2 and 3 (results are in Table 5).

The analysis described so far was performed on the water system under examination using a laptop having a processor of 1.74 GHz and a RAM of 4 GB and took a *Matlab* running time of roughly 24 seconds.

The recursive bisection procedure was now applied to the three large independent districts separately. The number *k* of *DMA*s to be created was fixed equal to 9, 4 and 3 for the first, the second and the third large independent districts respectively. A minimum required hydraulic head h_{req} of 20 m (30 pound per square inch) was considered. The solution obtained, which envisages the closure of 152 pipes for the creation of the boundaries of the 16 *DMA*s, is shown in Figure 13, while the size of the *DMA*s are reported in Table 6.

496

497 **CONCLUSIONS**

498 A new methodology for designing DMAs, based on graph theory, was proposed: it 499 uses graph theory principles and algorithms for automatically determining the boundaries of 500 each district. It allows the automatic creation of a feasible division of the network into the 501 required number of districts, from the start. At the same time, the constraints on the size 502 limits for a single district and the connectivity properties of the districts with the water source 503 are considered in the process of creating the DMAs. It ensures that the DMAs are roughly 504 equally sized in terms of water use, have an appropriate number of customer connections, and 505 are independent one from another, i.e., there are no connections between the districts, and 506 each one of them has a direct flow path to the water source. Finally, the performing of a 507 hydraulic simulation verifies the compliance of the minimum pressure requirements.

The methodology developed represents an improvement in respect to the ones found in literature: in fact each *DMA*, besides being characterized by an adequate number of customer connections, is also supplied directly by the transmission main and does not exchange water with adjacent *DMA*s. Therefore the distribution system is more reliable, the quality of the water delivered is higher and the possible spread of contaminants is greatly reduced.

514 Further improvements of the methodology could include node elevation among the 515 decision variables, in order to ensure that the difference in altitude of nodes belonging to the 516 same *DMA* is lower than an appropriate threshold. Secondly, it would be interesting to check 517 the fire flows of the resulting water distribution network, as it has been shown that the 518 division into districts could cause a reduction in fire flows. However, it should be taken into 519 account that a higher number of decision variables would imply a higher level of complexity 520 to deal with, which might result in an increased computational effort and in a reduction of 521 efficiency.

522 The methodology developed was applied to a real case study, in order to test its 523 applicability and effectiveness in real large urban water distribution systems. Results 524 highlighted how the methodology successfully provides for the division of the water 525 distribution network into a predetermined number of districts that are characterized by the 526 desired properties: appropriate size, connection with the main transmission system, hydraulic 527 independence from each other (i.e. no flow paths are available between two districts) and 528 hydraulic feasibility of the resulting network. Therefore, the methodology proposed has been 529 proved to represent a useful tool for DMAs design in large water distribution networks.

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- 617

618 FIGURES

619







Figure 2 - The methodology proposed



Figure 4 - The recursive bisection procedure.





L [node ID]	water use [lps]	W [lps]	
6394	0.03	0.03	
6392	0.06	0.09	
6393	0.30	0.39	
6395	0.07	0.46	
6391	1.99	2.45	
6396	0.26	2.71	
6398	0.02	2.73	
6390	0.82	3.55	
6399	0.09	3.64	77
1225	322		
(222	0.02	12.22	
1325			
6265	0.00	162.62	
5439	0.38	163.00	
5440	0.37	163.37	
6263	0.00	163.37	
5441	0.46	163.83	our!
6225	0.00	163.83	
7065	0.05	163.89	
6224	0.00	163.89	
6226	0.00	163.89	
7029	0.23	164.12	
7063	0.01	164.13	
1680	0.11	164.23	
7028	0.33	164.56	
7032	0.25	164.81	







Figure 9⁻ The water distribution network in analysis and its transmission main system.



 $\begin{array}{c} 656\\ 657\\ 658\end{array}$ Figure 10 - Groups of nodes whose water use is lower than the minimum value for a DMA $(W_d < W_{d,min}). \end{array}$



 $\begin{array}{ll} 662 \\ 663 \end{array} Figure 11 - Groups of nodes having total water use that is suitable for being considered a \\ DMA (W_{d,min} < W_d < W_{d,MAX}).. \end{array}$

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667 Figure 12 - Large independent districts: groups of nodes having total W_d greater than W_{d,MAX}

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674 TABLES

Number of Nodes	12523
Number of Links	14822
Number of Reservoirs	2
Number of Tanks	2
Number of Pumps	4
Total pipe length	1844.04 km
Total water demand	1512 l/s
Average water demand	0.121 l/s

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Table 1 - Water distribution network's characteristics

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Table 2 - Minimum and maximum size of a DMA

DMA index	Water demand [l/s]	Number of nodes	Number of connections with the transmission main
1	17.9	163	3
2	14.5	75	1
3	10.2	134	3
4	13.8	113	3
5	16.4	94	2
6	10.3	78	4
7	79.0	573	8
8	34.4	293	6
9	29.3	415	9
10	77.0	566	19
11	26.9	221	2
12	17.5	95	1
13	24.7	139	1
14	14.3	136	4
15	13.9	232	1
16	10.5	65	1
17	42.4	511	11
18	27.6	209	3
19	24.8	408	2
20	9.5	103	6

Table 3 - Characteristics of the existing DMAs

DMA index	Water demand [l/s]	Number of nodes	Number of connections with the transmission main
1	453.8	3921	73
2	201.2	1356	13
3	164.8	851	9
Tabla	1 Characteristic	a of the larg	a indonandant districts

Table 4 - Characteristics of the large independent distri	cts
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DMA index	\mathbf{k}_{\min}	k MAX
1	6	56
2	3	13
3	3	9

Table 5 - Minimum and maximum number of DMAs that can be created in the large independent district

DMA index	Number of Nodes	Water use [l/s]	Number of connections with the transmission mains
1	803	78.45	7
2	737	79.25	16
3	427	53.74	8
4	252	39.84	5
5	383	44.72	11
6	552	55.98	8
7	277	39.46	8
8	234	31.28	6
9	256	31.12	4
10	354	58.82	3
11	329	46.62	1
12	344	52.6	4
13	329	43.23	1
14	311	50.85	4
15	174	39.15	1
16	366	74.81	4

682Table 6 - Size of the DMAs created with the methodology and number of pipes connecting683each of them with the transmission mains