

# On the Two-Parabolic Subgroups of $SL(2, \mathbb{C})$

Sobre los subgrupos dos-parabólicos de  $SL(2, \mathbb{C})$

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ABSTRACT. We consider homomorphisms  $H_t$  from the free group  $F$  of rank 2 onto the subgroup of  $SL(2, \mathbb{C})$  that is generated by two parabolic matrices. Up to conjugation,  $H_t$  depends only on one complex parameter  $t$ . We study the possible relators, that is, the words  $w \in F$  with  $w \neq 1$  such that  $H_t(w) = I$  for some  $t \in \mathbb{C}$ .

We find several families of relators. Of particular interest here are relators connected with 2-bridge knots, which we consider in a purely algebraic setting. We describe an algorithm to determine whether a given word is a possible relator.

*Key words and phrases.* Representation, Parabolic, Wirtinger presentation, Two-generated groups, Homomorphism, Longitude.

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RESUMEN. Consideramos homomorfismos  $H_t$  del grupo libre  $F$  de rango 2 sobre el subgrupo de  $SL(2, \mathbb{C})$  que es generado por dos matrices parabólicas. Salvo conjugación,  $H_t$  depende sólo de un parámetro complejo  $t$ . Estudiamos los posibles relatores, esto es, las palabras  $w \in F$  con  $w \neq 1$  tal que  $H_t(w) = I$  para algún  $t \in \mathbb{C}$ .

Encontramos varias familias de relatores. De particular interés aquí son los relatores asociados con nudos de 2 puentes, los cuales consideramos de forma puramente algebraica. Describimos un algoritmo para determinar cuándo una palabra dada es un posible relator.

*Palabras y frases clave.* Representación, parabólico, presentación de Wirtinger, grupos dos-generados, homomorfismos, longitud.

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## 1. Introduction

The subgroup of  $\mathrm{SL}(2, \mathbb{C})$  generated by  $A_t = \begin{bmatrix} 1 & 0 \\ t & 1 \end{bmatrix}$  and  $B = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix}$  has been studied by many mathematicians, for instance by R. Riley [16, 17], J. Gilman [4] and P. Waterman [6]. It is of particular interest in knot theory [12, Chapter 4][2].

In terms of the corresponding Moebius transformations  $\alpha$  and  $\beta$  it is, up to conjugation, the only subgroup of  $\mathrm{PSL}(2, \mathbb{C})$  generated by two parabolic transformations with distinct fixed points. Indeed, we may assume that  $\alpha(0) = 0$ ,  $\beta(\infty) = \infty$ , moreover that  $\beta(z) = z + 1$ . Writing  $\alpha(z) = (az + b)/(cz + d)$  with  $ad - bc = 1$ , we may also assume that  $\mathrm{tr} \alpha = a + d = 2$ . Since  $b = 0$  it follows that  $a = d = 1$  and  $c = t \in \mathbb{C} \setminus \{0\}$  remains as a free parameter. It is convenient to allow  $t = 0$ .

Let  $F$  be the free group  $\langle x, y \rangle$ . We consider the homomorphisms  $H_t : F \rightarrow \mathrm{SL}(2, \mathbb{C})$  with  $H_t(x) = A_t$  and  $H_t(y) = B$ . For clarity we distinguish between the abstract group  $F$  and its image in the matrix group  $\mathrm{SL}(2, \mathbb{C})$ . Our main interest is to study the set of possible relators, that is the sets

$$R^\pm = \{r \in F, r \neq 1 \mid \text{there is } s \in \mathbb{C} \text{ with } H_s(r) = \pm I\}.$$

If  $r \in R^+$  then  $H_s(F)$  has a presentation  $\langle x, y; r_1, r_2, \dots \rangle$  with  $r_1 = r$  and perhaps other relators  $r_2, \dots$

We shall exhibit various families of relators, some old, some new. An important family of relators comes from the presentations  $\langle x, y; xw = wy \rangle$  of 2-bridge knots. Riley introduced the automorphism  $w \in F \rightarrow \tilde{w} \in F$  induced by  $x \rightarrow x^{-1}$  and  $y \rightarrow y^{-1}$ . Our group has the special property that  $r \in R^\pm$  implies  $\tilde{r} \in R^\pm$ . At the end we give 8 examples to illustrate the results and show their scope.

We try to give a systematic account of the theory including some folklore results. We will use several results from Combinatorial Group Theory [13, 10] and stress the connections to Knot Theory [2]. We have not been able to elucidate the role of palindromes, that is, words of  $F$  that read the same way forwards and backwards [8, 5].

We will not discuss the set orthogonal to  $R^+$ , namely

$$\{s \in \mathbb{C} \mid \text{there exists } r \neq 1 \text{ such that } H_s(r) = I\},$$

the set of  $s$  where  $H_s$  is not injective. This set and its closure has been studied in [16, 4, 6] and, in a more general context, in [18, 14].

## 2. Groups and Homomorphisms

Let  $SL(2, \mathbb{C})$  and  $PSL(2, \mathbb{C})$  be the groups with elements of the forms

$$C = \begin{bmatrix} a & b \\ c & d \end{bmatrix}, \quad \gamma(z) = \frac{az + b}{cz + d}, \quad (a, b, c, d \in \mathbb{C}, ad - bc = 1)$$

respectively. For  $t \in \mathbb{C}$  we write

$$A_t = \begin{bmatrix} 1 & 0 \\ t & 1 \end{bmatrix}, \quad B = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix}, \quad I = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}. \quad (1)$$

While the group  $PSL(2, \mathbb{C})$  of Moebius transformations is perhaps more important in analysis and geometry, the group  $SL(2, \mathbb{C})$  is more convenient for computations.

Let  $F$  be the abstract free group  $\langle x, y \rangle$ . There are [11, Theorem 8.04] unique homomorphisms

$$H_t : F \rightarrow SL(2, \mathbb{C}) \quad \text{with} \quad H_t(x) = A_t, H_t(y) = B. \quad (2)$$

$$h_t : F \rightarrow PSL(2, \mathbb{C}) \quad \text{with} \quad h_t(x) = \alpha_t, h_t(y) = \beta, \quad (3)$$

where  $\alpha_t(z) = z/(tz + 1)$  and  $\beta(z) = z + 1$  are both parabolic.

Every word  $w \neq 1$  in  $F$  can be uniquely written as

$$w = x^{e_0} y^{e_1} \dots x^{e_{m-1}} y^{e_m}, \quad e_\mu \in \mathbb{Z} \setminus \{0\} \quad (\mu = 1, \dots, m-1) \quad (4)$$

with  $e_0, e_m \in \mathbb{Z}$  and  $m \in \mathbb{N}$ . The exponent sums

$$\sigma_x(w) = e_0 + e_2 + \dots + e_{m-1}, \quad \sigma_y(w) = e_1 + e_3 + \dots + e_m \quad (5)$$

are invariant under conjugations in  $F$ . As in [17, p.206] we write  $\tilde{1} = 1$  and

$$\tilde{w} = x^{-e_0} y^{-e_1} \dots x^{-e_{m-1}} y^{-e_m}. \quad (6)$$

This defines an automorphism of  $F$ . In formulas we write  $(\cdot)^\sim$ .

Now suppose that  $w \in F$  is not conjugate in  $F$  to  $x^k$  or  $y^k$  with  $k \in \mathbb{Z}$ . Then the process of cyclic reduction shows that  $w$  is conjugate to a word  $u$  of the form

$$u = y^{k_1} x^{j_1} \dots y^{k_n} x^{j_n}, \quad k_\nu, j_\nu \in \mathbb{Z} \setminus \{0\} \quad (\nu = 1, \dots, n) \quad \text{with} \quad n \in \mathbb{N}. \quad (7)$$

Our standard form (7) allows us to arrange all words of  $F$  up to conjugation in sequences. Let  $(j_n)$  and  $(k_n)$  be any given sequences with  $j_n, k_n \in \mathbb{Z} \setminus \{0\}$ . Then we define  $u_0 = 1$  and

$$u_n = y^{k_1} x^{j_1} \dots y^{k_n} x^{j_n}, \quad \text{for} \quad n \in \mathbb{N}. \quad (8)$$

A conjugate of every  $w \in F$  will appear in many such sequences. No  $u_n$  with  $n \geq 1$  is conjugate to  $x^k$  or  $y^k$ .

Now we turn to the matrix elements. We often write

$$H_t(u) = \begin{bmatrix} a(t) & b(t) \\ c(t) & d(t) \end{bmatrix} \quad \text{or} \quad \begin{bmatrix} a & b \\ c & d \end{bmatrix}; \quad (9)$$

the elements of  $H_t(u_n)$  will be  $a_n, \dots, d_n$ .

The following proposition is folklore.

**Proposition 1.** *Let  $(u_n)$  be given by (8). Then, for  $n \geq 0$ ,*

$$a_{n+1} = j_{n+1}k_{n+1}ta_n + a_n + j_{n+1}tb_n, \quad (10)$$

$$b_{n+1} = k_{n+1}a_n + b_n, \quad (11)$$

$$c_{n+1} = j_{n+1}k_{n+1}tc_n + c_n + j_{n+1}td_n, \quad (12)$$

$$d_{n+1} = k_{n+1}c_n + d_n. \quad (13)$$

For  $n \geq 1$ , the  $a_n, \dots, d_n$  are polynomials over  $\mathbb{Z}$  of the forms

$$a_n = j_1k_1 \cdots j_nk_nt^n + \cdots, \quad b_n = j_1k_1 \cdots j_{n-1}k_{n-1}k_nt^{n-1} + \cdots, \quad (14)$$

$$c_n = j_1j_2k_2 \cdots j_nk_nt^n + \cdots, \quad d_n = j_1j_2k_2 \cdots j_{n-1}k_{n-1}k_nt^{n-1} + \cdots. \quad (15)$$

The trace  $\text{tr } H_t(u_n) = a_n + d_n$  is non-constant for  $n \geq 1$ .

*Proof.* By (8), (1) and (2) we have

$$\begin{aligned} H_t(u_{n+1}) &= H_t(u_n y^{k_{n+1}} x^{j_{n+1}}) = H_t(u_n) B^{k_{n+1}} A_t^{j_{n+1}} \\ &= \begin{bmatrix} a_n & b_n \\ c_n & d_n \end{bmatrix} \begin{bmatrix} 1 & k_{n+1} \\ 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ j_{n+1}t & 1 \end{bmatrix} \end{aligned}$$

and (10)–(13) follow by (9); we have  $a_0 = d_0 = 1$  and  $b_0 = c_0 = 0$ .

Now we prove the other assertions by induction. In each of the recursion formulas (10)–(13), all coefficients are in  $\mathbb{Z}$ . Furthermore, the degree of the first term is always, by induction hypothesis, higher than the degree of the other terms. Hence, by (14),  $a_{n+1}$  begins with  $j_{n+1}k_{n+1}t \cdot j_1k_1 \cdots j_nk_nt^n$  and  $b_{n+1}$  begins with  $k_{n+1} \cdot j_1k_1 \cdots j_nk_nt^n$ , similarly for  $c_{n+1}$  and  $d_{n+1}$ . The statement about the trace follows from (14) and (15).  $\square$

**Theorem 2.** *Let  $u_n$  satisfy (8) with  $|j_\nu| = |k_\nu| = 1$ . Then the coefficients  $a_{n,m}$  of  $a_n(t)$  and so on satisfy*

$$|a_{n,m}| \leq \binom{n+m}{2m} (0 \leq m \leq n), \quad |b_{n,m}| \leq \binom{n+m}{2m+1} (0 \leq m \leq n-1),$$

$$|c_{n,m}| \leq \binom{n+m-1}{2m-1} (1 \leq m \leq n), \quad |d_{n,m}| \leq \binom{n+m-1}{2m} (0 \leq m \leq n-1).$$

If  $u_n = (yx)^n$  then equality holds without taking absolute values.

Compare [16, p. 233] for the last statement. The values for the case  $v_n = (yx)^n$  were found by the method of generating functions. We have for instance

$$\sum_{n=0}^{\infty} a_n(t)z^n = \frac{1-z}{(1-z)^2 - tz}.$$

*Proof.* Let  $n \geq 0$  and  $m \geq 0$ . Writing  $a_{n,-1} = \dots = d_{n,-1} = 0$ , we obtain from (10)–(13) that

$$\begin{aligned} a_{n+1,m} &= a_{n,m} + j_{n+1}k_{n+1}a_{n,m-1} + j_{n+1}b_{n,m-1}, \\ b_{n+1,m} &= k_{n+1}a_{n,m} + b_{n,m}, \\ c_{n+1,m} &= c_{n,m} + j_{n+1}k_{n+1}c_{n,m-1} + j_{n+1}d_{n,m-1}, \\ d_{n+1,m} &= k_{n+1}c_{n,m} + d_{n,m}. \end{aligned}$$

Now we verify the assertions by induction on  $n$ . The case  $n = 0$  is clear because  $H_t(1) = I$ . We repeatedly use that  $\binom{\alpha}{\beta} + \binom{\alpha}{\beta-1} = \binom{\alpha+1}{\beta}$ .

Since  $|j_{n+1}| = |k_{n+1}| = 1$  the above recursion formulas show that

$$\begin{aligned} |a_{n+1,m}| &\leq |a_{n,m}| + |a_{n,m-1}| + |b_{n,m-1}| \\ &\leq \binom{n+m}{2m} + \binom{n+m-1}{2m-2} + \binom{n+m-1}{2m-1} \\ &= \binom{n+m}{2m} + \binom{n+m}{2m-1} = \binom{n+1+m}{2m}, \end{aligned}$$

$$|b_{n+1,m}| \leq |a_{n,m}| + |b_{n,m}| \leq \binom{n+m}{2m} + \binom{n+m}{2m+1} = \binom{n+1+m}{2m+1}$$

$$\begin{aligned} |c_{n+1,m}| &\leq |c_{n,m}| + |c_{n,m-1}| + |d_{n,m-1}| \\ &\leq \binom{n+m-1}{2m-1} + \binom{n+m-2}{2m-3} + \binom{n+m-1}{2m-2} = \binom{n+m}{2m-1}, \end{aligned}$$

$$|d_{n+1,m}| \leq |c_{n,m}| + |d_{n,m}| \leq \binom{n+m-1}{2m-1} + \binom{n+m-1}{2m} = \binom{n+m}{2m}.$$

If  $u_n = (yx)^n$  then  $j_{n+1} = k_{n+1} = 1$  and all quantities are non-negative. Hence we have equality in all the above inequalities.  $\square$

Our homomorphism has an important property with respect to the automorphism  $u \rightarrow \tilde{u}$  defined in (6). The following proposition is well-known.

**Proposition 3.** *Let  $u \in F$ ,  $H_t(u) = \begin{bmatrix} a & b \\ c & d \end{bmatrix}$  and  $Q = \begin{bmatrix} i & 0 \\ 0 & -i \end{bmatrix}$ . Then*

$$H_t(\tilde{u}) = QH_t(u)Q^{-1} = \begin{bmatrix} a & -b \\ -c & d \end{bmatrix}. \quad (16)$$

*Proof.* It is easy to see that

$$Q \begin{bmatrix} a & b \\ c & d \end{bmatrix} Q^{-1} = \begin{bmatrix} a & -b \\ -c & d \end{bmatrix}. \quad (17)$$

Hence it follows from (1) and (2) that

$$\begin{aligned} QA_tQ^{-1} &= \begin{bmatrix} 1 & 0 \\ -t & 1 \end{bmatrix} = A_t^{-1} = H_t(x^{-1}) = H_t(\tilde{x}), \\ QBQ^{-1} &= \begin{bmatrix} 1 & -1 \\ 0 & 1 \end{bmatrix} = B^{-1} = H_t(y^{-1}) = H_t(\tilde{y}). \end{aligned}$$

This implies (16) because  $(uv)\tilde{\phantom{u}} = \tilde{u}\tilde{v}$ ,  $(u^{-1})\tilde{\phantom{u}} = \tilde{u}^{-1}$  and  $H_t$  is a homomorphism.  $\square$

### 3. Relators

Our main interest is in the set of words

$$R^\pm = \{r \in F, r \neq 1 \mid \text{there exists } s \in \mathbb{C} \text{ with } H_s(r) = \pm I\}, \quad (18)$$

$$R = \{r \in F, r \neq 1 \mid \text{there exists } s \in \mathbb{C} \text{ with } h_s(r) = \text{id}\}. \quad (19)$$

Let  $N(r)$  denote the normal closure of  $r$ , that is the smallest normal subgroup of  $F$  with  $r \in N(r)$ .

It follows from Proposition 3 that  $H_s(u) = I$  implies  $H_s(\tilde{u}) = I$ ; see (6) for the definition of  $\tilde{u}$ . Hence, for  $r \in R^+$  or  $r \in R$ , the normal closure of  $\{r, \tilde{r}\}$  also belongs to  $R^+$  or  $R$  for the same  $s$ . Thus we have

$$u = v_1 r_1 v_1^{-1} \cdots v_m r_m v_m^{-1} \in R^+ \quad \text{or} \quad u \in R \quad (20)$$

where  $v_\mu \in F$  and  $r_\mu \in \{r, r^{-1}, \tilde{r}, \tilde{r}^{-1}\}$  for  $\mu = 1, \dots, m$  and  $m \in \mathbb{N}$ . It follows that the exponent sums (5) for  $u \in \mathbb{N}$  satisfy

$$\sigma_x(u) = \lambda \sigma_x(r), \quad \sigma_y(u) = \lambda \sigma_y(r) \quad \text{for some} \quad \lambda \in \mathbb{Z}. \quad (21)$$

If  $r \in R$  and thus  $h_s(r) = \text{id}$  for some  $s \in \mathbb{C}$  then, by the first isomorphism theorem, there is a homomorphism

$$h_{r,s} : \langle x, y; r \rangle \xrightarrow{\text{onto}} h_s(F) \quad (22)$$

defined by  $h_{r,s}(w) := h_s(u)$  for any  $w \in uN(r)$ . Now we show that  $h_{r,s}$  is in general not an isomorphism so that the representation of  $\langle x, y; r \rangle$  is not faithful. See Example 1.

**Proposition 4.** *If  $h_{r,s}$  is an isomorphism then  $\tilde{r}$  is conjugate to  $r$  or  $r^{-1}$ .*

Note that, if  $\tilde{r}$  is conjugate to  $r$ , then  $\sigma_x(r) = \sigma_y(r) = 0$ . This result is related to [3, Theorem 3.1]. On the other hand, it is easy to see that even  $\tilde{r} = r^{-1}$  holds if  $r$  is a palindrome, that is, the word  $r$  reads forwards the same as backwards. See [3, Propositions 3.4 and 3.2] for a fuller description.

*Proof.* Let  $h_{r,s}$  be an isomorphism. Since  $h_s(\tilde{r}) = \text{id}$  it follows that  $\tilde{r} \in N(r)$ . Hence the normal closure  $N(\tilde{r})$  of  $\tilde{r}$ , the smallest normal subset of  $F$  containing  $\tilde{r}$ , satisfies  $N(\tilde{r}) \subset N(r)$ . Furthermore  $\tilde{r} = v_1 r^{\pm 1} v_1^{-1} \dots v_m r^{\pm 1} v_m^{-1}$  and therefore

$$r = (\tilde{r})^{\sim} = \tilde{v}_1 \tilde{r}^{\pm 1} \tilde{v}_1^{-1} \dots \tilde{v}_m \tilde{r}^{\pm 1} \tilde{v}_m^{-1} \in N(\tilde{r}).$$

Hence  $N(r) \subset N(\tilde{r})$  so that  $r$  and  $\tilde{r}$  have the same normal closure  $N(r)$ . It follows [13, p. 261] [10, Proposition. 5.8, p. 106] that  $\tilde{r}$  is conjugate to  $r$  or  $r^{-1}$ .  $\square$

Now we present a general method to obtain relators. See Examples 2 and 3.

**Theorem 5.** *Let  $u \in F$  not be conjugate to  $x^k$  or  $y^k$  with  $k \in \mathbb{Z}$ . Then*

$$u\tilde{u} \in R^+, \quad u^2 \in R^-, \tag{23}$$

$$u^n \in R^+ \cap R^-, \quad \text{for } n \geq 3, \tag{24}$$

thus  $u\tilde{u} \in R$  and  $u^n \in R$  for  $n \geq 2$ .

For  $u\tilde{u}$  we have to exclude the case that  $u$  is a palindrome because then  $u\tilde{u} = 1$ , compare (18) and (19). Note that  $u\tilde{u} = 1$  holds if and only if  $u$  is a palindrome.

*Proof.*

(a) We obtain from (9) and (16) that

$$H_t(u\tilde{u}) = \begin{bmatrix} a & b \\ c & d \end{bmatrix} \begin{bmatrix} a & -b \\ -c & d \end{bmatrix} = \begin{bmatrix} 1 + a(a-d) & -b(a-d) \\ c(a-d) & 1 - d(a-d) \end{bmatrix}.$$

Since  $a-d$  is non-constant by Proposition 1, it follows that  $a(s) - d(s) = 0$  for some  $s \in \mathbb{C}$ . Hence  $H_s(u\tilde{u}) = I$ .

(b) First let  $n \geq 2$ . There exists  $s$  such that  $\text{tr } H_s(u) = 2 \cos(\pi/n)$ . Then  $\text{tr } H_s(u) \neq \pm 2$ . Hence  $H_s(u)$  is conjugate to  $\text{diag}(e^{i\pi/n}, e^{-i\pi/n}) \neq I$ . It follows that  $H_s(u^n) = -I$ .

Now let  $n \geq 3$ . There exists  $s$  such that  $\text{tr } H_s(u) = 2 \cos(2\pi/n)$  so that, again,  $\text{tr } H_s(u) \neq \pm 2$ . Hence  $H_s(u)$  is conjugate to  $\text{diag}(e^{2\pi i/n}, e^{-2\pi i/n}) \neq I$  so that  $H_s(u^n) = I$ .  $\square$

**Proposition 6.** *Let  $u \in F$  have the form (7). If  $s \neq 0$  and  $H_s(u) = I$  then  $s$  is an algebraic number of degree  $\leq (n-1)/2$ , and if  $H_s(u) = -I$  then  $s$  is an algebraic number of degree  $\leq n/2$ . If  $|j_\nu| = |k_\nu| = 1$  for  $\nu = 1, \dots, n$  then  $s$  is an algebraic integer.*

This is common knowledge except for the sharp bounds  $(n-1)/2$  and  $n/2$  for the degrees, see Example 4. Note that  $H_s(F) \subset \mathrm{SL}(2, \mathbb{Z}[s])$ . We have shown that, up to conjugation,  $u$  has the form (7) whenever  $u$  is not conjugate to  $x^k$  or  $y^k$ .

*Proof.* We use the notation (9). If  $H_s(u) = \pm I$  then  $c(s) = 0$ . Hence, by (15),  $s$  is an algebraic number which is an algebraic integer if  $|j_\nu| = |k_\nu| = 1$ . Now let  $p(t)$  be the minimal polynomial of  $s$ . First let  $H_s(u) = I$  and  $s \neq 0$ . We write

$$a + d - 2 = -(a-1)(d-1) + bc.$$

Since  $a(s) = d(s) = 1$  and  $b(s) = c(s) = 0$ , we conclude that  $p$  divides  $(a-1)$ ,  $(d-1)$ ,  $b$  and  $c$ . Hence  $p^2$  divides  $a + d - 2$ . Since  $s \neq 0$ , it follows that  $tp^2$  divides  $a + d - 2$ , which is a polynomial of degree  $n$  by (14). Thus  $s$  has degree  $\leq (n-1)/2$ .

Now let  $H_s(u) = -I$ . We write

$$a + d + 2 = (a+1)(d+1) - bc.$$

Now  $p$  divides  $(a+1)$ ,  $(d+1)$ ,  $b$  and  $c$ . Hence  $p^2$  divides the polynomial  $a + d + 2$  of degree  $n$ . Thus  $s$  has degree  $\leq n/2$ .  $\square$

Now we describe an *algorithm to determine* whether  $u \in F$  belongs to  $R^+$  or  $R^-$ . This is not the case if  $u$  is conjugate to  $x^k$  or  $y^k$ . Therefore we may assume that  $u$  has the form (7). We use the notation (9).

First we check whether it is possible that  $b = c = 0$  for some  $t \in \mathbb{C}$ . To do this we calculate the polynomial

$$q_0 := \gcd(b, c) \in \mathbb{Z}[t]. \quad (25)$$

If  $\deg q_0 = 0$  then  $u \notin R^+ \cup R^-$ . If however  $\deg q_0 > 0$  then we calculate the polynomials

$$q^\pm := \gcd(a \mp 1, q_0). \quad (26)$$

If  $\deg q^+ = 0$  then  $u \notin R^+$ , if  $\deg q^- = 0$  then  $u \notin R^-$ .

Now if  $\deg q^\pm > 0$  then there is  $s \in \mathbb{C}$  such that  $a(s) = \pm 1$ . It follows from (25) and (26) that  $b(s) = c(s) = 0$  so that  $1 = a(s)d(s) - b(s)c(s) = \pm d(s)$ . Therefore we have  $H_s(u) = \pm I$  and thus  $u \in R^\pm$ . Additionally we may factorize  $q^\pm$  into irreducible polynomials over  $\mathbb{Z}$ . If  $s$  is a zero of a factor then all other zeros  $t$  of this factor satisfy  $H_t(u) = \pm I$ . The main computational difficulty of this algorithm is that very large integer coefficients may occur during the calculation of (25) and (26).

#### 4. The Wirtinger Relators and the Longitude

Let  $K$  be a knot in  $\mathbb{R}^3$ , see e.g. [2, 9]. The complement  $\Omega = \overline{\mathbb{R}^3} \setminus \overline{V(K)}$ , where  $V(K)$  is a tubular neighborhood of  $K$ , is a multiply connected domain. The fundamental group  $\Pi_1(\Omega)$  is an important invariant of  $K$  though it does not completely determine the equivalence class of  $K$ , although the prime knots are determined by their knot group [7]. A very well understood family of knots are the so called 2-bridge knots and links [20, 2, 15, 19].

The fundamental group of a 2-bridge knot admits a presentation  $\langle x, y; xw_n = w_n y \rangle$  where

$$\begin{aligned} w_n &= y^{k_n} x^{k_{n-1}} \dots y^{k_1} x^{k_1} y^{k_2} \dots y^{k_{n-1}} x^{k_n}, & k_\nu &\in \{1, -1\}, & n \text{ odd} \\ w_n &= y^{k_n} x^{k_{n-1}} \dots y^{k_2} x^{k_1} y^{k_1} \dots y^{k_{n-1}} x^{k_n}, & k_\nu &\in \{1, -1\}, & n \text{ even} \end{aligned} \quad (27)$$

for  $n \in \mathbb{N}$ , where  $k_\nu, \nu = 1, \dots, n$ , satisfy some additional conditions [2, 1]. We inverted the usual order of exponents in order to have a recursive definition. On the following we leave the context of knot theory and call any word of the form (27) a *Wirtinger word*.

It follows from (27) that, with  $\tilde{\phantom{x}}$  defined in (6),

$$w_{n+1} = (y^{-k_{n+1}} \tilde{w}_n x^{-k_{n+1}}) \tilde{\phantom{x}} \quad (28)$$

Instead of (9) we now write

$$W_n = H_t(w_n) = \begin{bmatrix} a_n & b_n \\ c_n & d_n \end{bmatrix}. \quad (29)$$

It follows from (28), (1) and Proposition 3 that, with  $k = k_{n+1}$ ,

$$\begin{aligned} W_{n+1} &= QB^{-k} W_n^{-1} A^{-k} Q^{-1} \\ &= Q \begin{bmatrix} 1 & -k \\ 0 & 1 \end{bmatrix} \begin{bmatrix} d_n & -b_n \\ -c_n & a_n \end{bmatrix} \begin{bmatrix} 1 & 0 \\ -kt & 1 \end{bmatrix} Q^{-1}. \end{aligned}$$

Using also (17) and  $k^2 = 1$ , we obtain

$$W_{n+1} = \begin{bmatrix} ta_n + ktb_n + kc_n + d_n & ka_n + b_n \\ kta_n + c_n & a_n \end{bmatrix}. \quad (30)$$

Since  $b_0 = c_0 = 0$  we deduce by induction the well-known formula [12, p. 141]

$$c_n = tb_n. \quad (31)$$

Hence we obtain from (30) the recursion formulas

$$a_{n+1} = ta_n + 2k_{n+1}tb_n + a_{n-1}, \quad b_{n+1} = k_{n+1}a_n + b_n.$$

Now (27) is a special case of (7). Hence the estimates of Theorem 2 apply also with the new notation.

Now we drop the index  $n$  and write

$$W_t := H_t(w) = \begin{bmatrix} a(t) & b(t) \\ c(t) & d(t) \end{bmatrix}. \quad (32)$$

**Theorem 7.** *If  $w$  satisfies the Wirtinger condition (27) then*

$$H_s(xw) = H_s(\tilde{w}y^{-1}) \quad (33)$$

*holds if and only if  $a(s) + 2b(s) = 0$ . Thus  $wy\tilde{w}^{-1}x \in R^+$ .*

*Proof.* By (16) and (31), the condition (33) is equivalent to

$$\begin{bmatrix} a & b \\ t(a+b) & tb+d \end{bmatrix} = A_t W_t = \tilde{W}_t B^{-1} = \begin{bmatrix} a & -a-b \\ -tb & tb+d \end{bmatrix},$$

and this condition holds if and only if  $t$  satisfies  $a(t) + 2b(t) = 0$ . The non-constant polynomial  $a + 2b$  has a root  $s$ . Hence

$$r = wy\tilde{w}^{-1}x = x^{-1}(xw)(\tilde{w}y^{-1})^{-1}x \in R^+. \quad \checkmark$$

In knot theory the condition (33) is replaced by

$$H_s(xw) = H_s(wy), \quad r = wy^{-1}w^{-1}x \in R^+. \quad (34)$$

which holds if and only if  $a(s) = 0$ , see e.g. [12, p. 141].

By Proposition 3 and conjugation, we see that (34) implies

$$H_s(x\tilde{w}) = H_s(\tilde{w}y), \quad \tilde{r} = \tilde{w}y\tilde{w}^{-1}x^{-1} \in R^+. \quad (35)$$

Condition (34) induces a homomorphism  $h_{r,s}$  from  $\langle x, y; r \rangle$  into  $\text{PSL}(2, \mathbb{C})$ ; see (22). Now (35) says that it automatically induces a homomorphism from  $\langle x, y; r, \tilde{r} \rangle$ . In the case of a 2-bridge knot it is known [2, 1] that there exists a faithful discrete  $\text{SL}(2, \mathbb{C})$ -representation of a 2-bridge knot of type  $(p, q)$  with  $q \not\equiv \pm 1$ , so that  $xw = wy$  implies  $x\tilde{w} = \tilde{w}y$ . But this is not true in general, see Example 5.

The situation is different for (33) because  $\tilde{r} = \tilde{w}y^{-1}w^{-1}x^{-1}$  is conjugate to  $r^{-1}$  so that  $r$  and  $\tilde{r}$  have the same normal closure; compare Proposition 4. Hence they induce the same group.

For 2-bridge knots the group  $G = \langle x, y; xw = wy \rangle$  and its peripheral subgroup are important concepts to distinguish equivalence classes of knots. This subgroup is generated by a meridian, say  $y$ , and the *longitude*  $l = w^{-1}\tilde{w}$

(see [17, p. 206]). We omitted Riley's factor  $y^{2\sigma}$ . It is easy to check that  $(r = 1, \tilde{r} = 1)$  is equivalent to  $(r = 1, ly = yl)$ .

Now we study the longitude  $l = w^{-1}\tilde{w}$  in a more general context. We do not assume that the word  $w$  comes from knot theory and we do not assume the consequence (31) of the Wirtinger condition. For  $w \in F$  we obtain from (32) and (16) that

$$H_t(l) = W_t^{-1}\widetilde{W}_t = \begin{bmatrix} ad + bc & -2bd \\ -2ac & ad + bc \end{bmatrix}. \quad (36)$$

We note that  $l = w^{-1}\tilde{w}$  is a palindrome.

**Theorem 8.** *Let  $w$  satisfy (7) with  $|j_\nu| = |k_\nu| = 1$  and let  $a(s) = 0$ . Then*

$$L_s := H_s(l) = \begin{bmatrix} -1 & -2b(s)d(s) \\ 0 & -1 \end{bmatrix}. \quad (37)$$

*If  $a = c + d$  then  $b(s)d(s) = 1$ . If the polynomial  $a$  is irreducible and if  $a \neq c + d$  then  $b(s)d(s) \notin \mathbb{Q}$  and  $L_s$  and  $B$  generate a free abelian group of rank 2.*

Formulas similar to (37) follow from (36) if  $b(s) = 0$ ,  $c(s) = 0$  or  $d(s) = 0$ . The 2-bridge knots of type  $(2n + 1, 1)$  have the Wirtinger word  $w = (yx)^n$ . It follows from Theorem 2 that  $a = c + d$  so that  $b(s)d(s) = 1$ . See Examples 6, 7 and 8.

*Proof.* Since  $ad - bc = 1$  we can write  $ad + bc = -1 + 2ad$ . Hence (37) follows from (36). If  $a = c + d$  then  $c(s) = -d(s)$  and thus  $b(s)d(s) = -b(s)c(s) = 1$  because  $a(s) = 0$ .

Now let  $q := b(s)d(s)$  and suppose that  $q \in \mathbb{Q}$ . Since  $a(s) = 0$ , it follows from Proposition 1 that  $s$  is an algebraic integer so that  $q \in \mathbb{Z}$ . It follows from (14) and (15) that

$$f(t) := qc(t) + d(t) = q\lambda t^n + \dots, \quad \lambda = \pm 1. \quad (38)$$

Since  $a(s) = 0$  implies  $b(s)c(s) = -1$  we have

$$b(s)f(s) = qb(s)c(s) + b(s)d(s) = -q + q = 0$$

so that  $f(s) = 0$ . Hence the irreducible polynomial  $a(t)$  divides  $f(t)$ . Since  $a(t) = \lambda t^n + \dots$  with the same  $\lambda$ , we conclude from (38) that  $q = 1$  and therefore  $a = c + d$ . If  $a \neq c + d$  we therefore have  $-2b(s)d(s) \notin \mathbb{Q}$  so that  $L_s$  and  $B$  are free abelian generators.  $\square$

### 5. Examples

The words of  $F$  in the following examples are generated by

$$z_0 = yx, \quad z_1 = yx^{-1}, \quad z_2 = y^{-1}x, \quad z_3 = y^{-1}x^{-1}.$$

All polynomials will be written as the product of irreducible factors in  $\mathbb{Z}[t]$ . The factorization used the program `Kash3` developed by M. Pohst and his group, [www.math.tu-berlin.de/~kant](http://www.math.tu-berlin.de/~kant).

**Example 1.** The following two words

$$\begin{aligned} r_1 &= z_0^2 z_1 z_3^2, & \sigma_x(r_1) &= -1, & \sigma_y(r_1) &= 1, \\ r_2 &= z_0^{10}, & \sigma_x(r_2) &= 10, & \sigma_y(r_2) &= 10 \end{aligned}$$

are relators with the same minimal polynomial  $1 + 3t + t^2$ . The normal closures satisfy  $r_1 \notin N(r_2)$  and  $r_2 \notin N(r_1)$  because the exponent sums do not satisfy (21). It follows that no homomorphism  $h_{r,s}$  with  $s = -1/2 \pm \sqrt{5}/2$  can be injective, see (22).

**Example 2.** Let  $u = z_0^2 z_2$  and  $r = u\tilde{u} = z_0^2 z_2 z_3^2 z_1$ . The polynomials for  $u$  are  $a(t) = 1 + 4t - t^2 - t^3$  and  $d(t) = 1 - t - t^2$ . Now part (a) of the proof of Theorem 5 shows that  $H_t(r) = I$  if and only if  $a(s) - d(s) = s(5 - s^2) = 0$ . Hence  $r \in R^+$ .

**Example 3.** Let  $r = z_0^2$ . Then

$$H_t(r) = \begin{bmatrix} 1 + 3t + t^2 & 2 + t \\ 2t + t^2 & 1 + t \end{bmatrix} \neq \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}, \quad \text{for } t \in \mathbb{C}$$

so that  $r \notin R^+$ . But (23) shows that  $r \in R^-$ .

**Example 4.** The following words belong to  $R^+ \cap R^-$ . Their minimal polynomials

$$\begin{aligned} u = z_0^5 : & \quad p^+(t) = 5 + 5t + t^2, & p^-(t) &= 1 + 3t + t^2, \\ u = z_0^6 : & \quad p^+(t) = 3 + 4t + t^2, & p^-(t) &= 2 + 9t + 6t^2 + t^3 \end{aligned}$$

have the smallest degrees possible by Proposition 6.

**Example 5.** The Wirtinger word  $w = z_0 z_1 z_1 z_0$  does not come from a 2-bridge knot. Its relator is  $r = wy^{-1}w^{-1}x$  with  $\sigma_x(r) = 1$ ,  $\sigma_y(r) = -1$ . Furthermore  $\tilde{r} = \tilde{w}y\tilde{w}^{-1}x^{-1}$  with  $\sigma_x(\tilde{r}) = -1$ ,  $\sigma_y(\tilde{r}) = 1$  so that  $\tilde{r}$  is not conjugate to  $r$ . Now  $r^{-1}$  contains  $y^{-1}x^{-1}y^{-1}x^{-1}$  whereas no conjugate of  $\tilde{r}$  contains this word. Hence  $r$  is not conjugate to  $r^{-1}$  either. Thus it follows from Proposition 4 that, with  $H_s(r) = I$ , the homomorphism

$$\langle x, y; r, \tilde{r} \rangle = \langle x, y; xw = wy, x\tilde{w} = \tilde{w}y \rangle \rightarrow \text{SL}(2, \mathbb{C})$$

is not injective.

**Example 6.** The Wirtinger word of the 2-bridge knot of type  $(9, 1)$  is  $w = z_0^4$  and

$$a(t) = (1+t)(1+9t+6t^2+t^3)$$

is reducible. It satisfies  $a = c + d$  and thus  $b(s)d(s) = 1$  by Theorem 8.

**Example 7.** Let  $w = z_0 z_3 z_2 z_0$ . This is not a Wirtinger word because  $c \neq tb$ . It satisfies

$$\begin{aligned} a(t) &= (1+t)p(t), & p(t) &= -1+t+2t^2+t^3, \\ b(t)d(t) - 1 &= (1+t)(-1-t-t^2+2t^3+2t^4+t^5), \\ b(t)d(t) + 1 &= (-1+t+t^2+t^3)p(t). \end{aligned}$$

Hence  $b(-1)d(-1) = 1$  whereas  $b(s)d(s) = -1$  if  $p(s) = 0$ . Thus, in Theorem 8, the assumption that  $a(t)$  is irreducible can not be omitted.

**Example 8.** The 2-bridge knot of type  $(5, 3)$  has  $w = z_1 z_2$  and  $a(t) = 1-t+t^2$ . This gives  $s = (1+i\sqrt{3})/2$  and  $b(s)d(s) = \pm i\sqrt{3}$ .

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