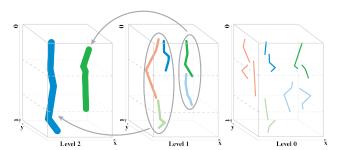
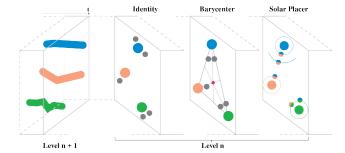
# A Multilevel Approach for Event-Based Dynamic Graph Drawing

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(a) Example of the coarsening process. When two or more nodes are merged together, their presence intervals are merged as well.

(b) Placement strategies for the three MultiDynNoS variants.

Figure 1: Multilevel strategies have two important stages: coarsening and placement. In this event-based multilevel approach, we coarsen and place trajectories. An example of the coarsening (a) and placement (b) stages used by the approach.

## Abstract

The timeslice is the predominant method for drawing and visualizing dynamic graphs. However, when nodes and edges have real coordinates along the time axis, it becomes difficult to organize them into discrete timeslices, without a loss of temporal information due to projection. Event-based dynamic graph drawing rejects the notion of a timeslice and allows each node and edge to have its own real-valued time coordinate. Nodes are represented as trajectories of adaptive complexity that are drawn directly in the three-dimensional space-time cube (2D + t). Existing work has demonstrated clear advantages for this approach, but these advantages come at a running time cost. In response to this scalability issue, we present MultiDynNoS, the first multilevel approach for event-based dynamic graph drawing. We consider three operators for coarsening and placement, inspired by Walshaw, GRIP, and FM<sup>3</sup>, which we couple with an event-based graph drawing algorithm. We evaluate our approach on a selection of real graphs, showing that it outperforms timeslice-based and existing event-based techniques.

### CCS Concents

• Human-centered computing → Graph drawings; Visualization;

## 1. Introduction

Usually, a dynamic graph is defined as a succession of individual static graphs [BBDW17], each one representing the state of the graph at a specific time instant (also known as a *timeslice*). This definition has two advantages: it works well for clearly defined time intervals (e.g., yearly, monthly, etc.) and allows for existing static layout algorithms to be used directly for drawing. However, when nodes and edges have independent time coordinates, projecting onto the nearest timeslice results in a quantization er-

ror, potentially reducing drawing quality (see this video). Event-based networks (also known as temporal networks [HS12]) do not suffer from this issue as real-valued time coordinates are specified for each node and edge. Unlike timeslice-based approaches, event-based drawing algorithms exploit the full temporal resolution of the data by optimizing node trajectories in the space-time cube (2D + t), outperforming timeslice-based techniques in terms of drawing quality [SAK17, SAK20] albeit with significant costs in terms of running time and computational resources. These higher running times have limited the use of event-based graph drawing on net-

works with a small number of events, despite the quality improvements over timeslice-based techniques.

This paper presents MultiDynNoS: the first multilevel event-based graph drawing algorithm, capable of bringing the time to draw event-based networks comparable to timeslice-based approaches. Similar to standard multilevel techniques for static graphs, MultiDynNoS follows a *coarsening-refinement* strategy. We adapt the coarsening and placement strategies of Walshaw [Wal03], GRIP [GK00], and FM³ [HJ04], designed for static graphs, to operate on node trajectories for drawing temporal graphs in the space-time cube. Our experiments show that MultiDyn-NoS drawing quality, in terms of stress, is comparable to existing event-based techniques [SAK17, SAK20] but with significant running time improvements, making them more competitive with timeslice-based approaches [BM11].

#### 2. Related Work

The visualization of dynamic graphs has been studied extensively [BBDW17] with animated techniques [APP10, AP16, FQ11, BPF14] and timeline approaches [APP10,SA06,BVB\*11,LHS\*15, AB20] receiving considerable attention. We focus on the closest related work to our work in this section.

**Multilevel Graph Drawing.** In the 2000s, multilevel graph drawing algorithms [Wal03, AMA07, GK00, HJ04, BGKM10] were devised to scale to larger static graphs. These algorithms construct a hierarchy of coarse graphs and exploit this hierarchy to accelerate the drawing. Multilevel graph drawing approaches have been adapted to an online dynamic setting [CCM17, Vel07, Cra16]. Multi-*layer* networks, where several node and edge layers have different meaning [MGM\*19], have been used for visualization.

Temporal Networks and Event-Based Visualization. Temporal and event-based networks [HS12, LVM18] have been studied extensively for automatic graph analysis. For most of the past two decades, visualization of temporal networks has focused on drawing a series of projected timeslices in a way that encourages a stable drawing [BBDW17] - the position of nodes and edges should change as little as possible when a change is made to the graph [CP96] so that nodes and edges can be easily identified [AP12, AP16]. Algorithms have been explored to optimize the simultaneous drawing of timeslices in offline [DG02, DGK01, EHK\*03, BM11] and online [MELS95, GDBG12, FT08] scenarios. Event-based visualization techniques [DSP\*17, MLMdO\*13, MLL\*131 visualise event sequences with real time coordinates for each data point. Event-based dynamic graph drawing algorithms directly draw these event-based/temporal graphs in the space-time cube [SAK17, SAK20]. Other techniques, such as HOTVis [PS21], exploit the temporal ordering of the edges (the causal paths) to influence the layout. However, they focus on 2D visualizations and do not optimize the drawing across the space-time cube.

**Contribution.** The literature indicates a growing interest in event-based visualizations of networks for visual analytics applications. Event-based dynamic graph drawings can potentially yield improved drawing quality over timeslice-based approaches, motivating our research on more scalable techniques for embedding temporal networks in the space-time cube.

### 3. MultiDynNoS Pipeline

Consider a temporal network D=(V,E) where each node and edge possesses a number of *attributes* which are functions of time. The *appearance* of a node  $v \in V$  is defined as  $A_v : V \times T \to [true, false]$  (edge appearance is defined similarly) which maps to node/edge insertion and deletion in the event-based graphs.  $A_v$  defines a series of intervals in T (time) in which the node/edge is present. The *position* of a node in the plane over time is defined as  $P_v : V \times T \to R^2$ . When defined in this way, the appearance and position of the nodes are represented as a series of trajectories through time embedded in the space-time cube (e.g., Fig. 1a): lines that define node movement in the two dimensional plane as time passes downwards in the cube. We also define a *flattened* graph as the weighted static counterpart of a temporal graph where node and edge weights represent the cumulative duration of the time intervals in which their appearance attribute function yields true.

**Layout Process.** First, a *coarsening* operator is applied on *D* to generate a *coarse hierarchy* of the graph. Starting from the coarsest graph, each level gets *refined*: its drawing is computed and its coordinates are used to *place* (i.e. assign the initial coordinates) the vertices on the level below. This initial placement provides quicker convergence in the next refinement cycle. Refinement ends when the layout for the input graph is computed.

Coarsening. Coarsening yields a hierarchy of coarse eventbased graphs  $D_H = \{D_w, D_1, ..., D_k\}$ , with "depth" k, to be used in the refinement stage.  $D_w$  is the input graph with an added attribute constant function representing the node and edge weights from the flattened graph  $D_f$ . The finest level is  $D_1$  and the coarsest  $D_k$ . For each level  $D_n = (V_n, E_n)$ , we order the vertices of  $V_n$  by their weight and put them on a stack. We pop the stack and get the heaviest vertex  $v_n$ : its copy  $v_{n+1}$  is then assigned to  $V_{n+1}$ . At this point, we select a subset of the neighbors of  $v_n$ , depending on the coarsening strategy, summing their weights and merging their appearance intervals with  $v_{n+1}$ . We refer to  $v_n$  as the "representative" in  $V_{n+1}$  of the vertices merged with it in  $V_n$ . We refer to the set of representatives of level n as  $\overline{V}_n$ . Once complete,  $v_n$  and the vertices merged with it are removed from the stack. This process is repeated until the stack is empty. Coarsening stops at the coarsest hierarchy level  $D_k$ when the node count falls below a threshold or it is  $\geq 95\%$  the size of level  $D_{k-1}$ . The latter condition is introduced to avoid a deep hierarchy with very similar level sizes, which would slow down drawing significantly. We implemented three different coarsening strategies. First, we implemented the Maximal Matching, found in the multilevel approach by Walshaw [Wal03], where pairs of vertices connected by an edge belonging to the graph maximal matching are merged together in each level. Second, we implemented the Maximal Independent Set coarsening, used by GRIP [GK00]. Once a vertex is selected to be part of the new level, it is merged together with all of its neighbors. Finally, we implemented the Solar Merger algorithm, used by FM<sup>3</sup> [HJ04]. Each selected vertex is merged together with its neighbors up to distance 2, creating a "Solar System Partitioning" of the graph. Once the vertex set for the new level is created, we generate  $E_{n+1}$ : for each edge  $e_n = (v_n, w_n)$ , we create an edge  $e_{n+1} = (v_{n+1}, w_{n+1})$  such that  $v_n$  and  $w_n$  were merged in  $v_{n+1}$  and  $w_{n+1}$  respectively. If that edge already exists, its presence is merged with the one of  $e_n$ .

Coarsest Level Placement. Initial placement assigns the initial coordinates of the trajectories in  $D_k$  as follows: we flatten  $D_k$  to obtain  $D_k'$  and draw it using a static layout algorithm. Node trajectories are centered in these newly computed coordinates and extruded vertically downwards across time. Subsequent refinement steps with an event-based drawing algorithm [SAK20] "bend" these trajectories to change the position of the nodes across time. A good initial placement is expected to yield smoother trajectories with few bends, which resolves in nodes with smoother movement.

**Refinement.** During each refinement iteration, Dyn-NoSlice [SAK20] is run on  $D_n$ . One of the key points of the multilevel strategy is that more quality-oriented layout parameters can be used on coarse graphs, since they are smaller in size and therefore quicker to draw. As the size of the graph to layout increases, speed can be emphasized. In our approach, we tune two parameters: the maximum node mobility and the number of layout algorithm iterations. Coarser levels will benefit from more flexible trajectories, while finer levels are more conservative with reduced iterations and movement. The parameters decrease linearly by 7% at each level. This value was obtained empirically when the considering quality/running time trade off. Time trajectory postprocessing of DynNoSlice [SAK17, SAK20] runs once every two layout iterations in the coarser levels and the interval grows by 2 with each new level. Once the layout for  $D_n$  is computed (and  $D_n \neq D_w$ ), the final coordinates are used to *place* the node trajectories on level  $D_{n-1}$ . First, each representative  $v_{n-1} \in \overline{V}_{n-1}$ is placed at the coordinates of the corresponding vertex in  $V_n$ . We compute the initial coordinates of the remaining vertices based on the new coordinates of their representative . We implemented three placement operators (Fig. 1b) inspired by Walshaw [Wal03], GRIP [GK00], and  $FM^3$  [HJ04]. The first strategy is the *identity* placer: the nodes are placed in the same position as their representative. The second strategy places the trajectories close to the barycenter of the coordinates of the representative's neighbors at level n+1. The final position of the node is skewed towards its own representative by a fixed rate. The third strategy is similar to barycenter but changes the attraction of the representative cluster. Specifically, given any two neighboring nodes  $v_{n+1}, w_{n+1} \in V_{n+1}$ , the solar system partitioning guarantees that representatives at level n,  $v_n$  and  $w_n$ , are at most distance 5 from each other. Since  $v_n$  and  $w_n$  neighbors up to distance 2 are merged together in the FM<sup>3</sup> coarsening, with this information it is possible to reconstruct the relative position of any of the merged trajectories in the paths between  $v_n$  and  $w_n$ , and place them accordingly. When the path position is not known it uses the barycenter placement strategy. For all approaches, randomness is added to the final coordinates to avoid possible accidental coordinate overlaps.

## 4. Experimental Evaluation

We conduct an evaluation where we repeat the experiment performed in DynNoSlice [SAK20] to compare MultiDynNoS to state-of-the-art dynamic graph layout algorithms on known metrics. Our research question can be formulated as follows: "Is MultiDynNoS faster than DynNoSlice, while providing layouts with comparable drawing quality?". We implemented MultiDynNoS and tested it (results in Table 1).

Metrics and Strategies. We evaluate the layouts using quality and readability metrics. We include: (i) the time, drawing time in seconds; (ii) Movement, the average distance travelled by a node during graph evolution [BM11, SAK20]; (iii) Crowding: the number of times nodes pass close to each other in the animation of the dynamic graph [SAK20]; (iv) Depth: coarsening depth (multilevel strategies only); (v) StressOn and (vi) StressOff, which are the layout stress computed on a per-timesliced basis or between timeslices, respectively, with optimal scaling [SAK20] applied. We test three MultiDynNoS variants: MultiDynNoS wi\_id is the Walshaw variant of MultiDynNoS with maximal matching of trajectories and identity placement; MultiDynNoS is\_gr is the GRIP variant of MultiDynNoS with maximal independent set coarsening of trajectories and barycenter placement; Multi-DynNoS sm\_sp is the FM<sup>3</sup> variant of MultiDynNoS with the FM<sup>3</sup> coarsening and placement strategy. Each variant is tested alternating the drawing algorithm for the coarsest level placement between sfdp [Hu05] and fdp [FR91]. The variants of MultiDynNoS are tested against Visone [BW04], a state-ofthe-art timeslice-based dynamic graph drawing algorithm, and DynNoSlice [SAK17, SAK20]. sfdp flat flattens the entire event-based data and draws it once as a static graph using

**Results.** Table 1 shows the results of our experiments. In terms of running time, on all the experiment instances MultiDynNoS is competitive with Visone and can be an order of magnitude faster than DynNoSlice. This represents a leap forward than previous studies [SAK17,SAK20] (whose results have been replicated here), where Visone always had the best performance when compared to DynNoSlice on this same set of graphs. In terms of drawing quality, MultiDynNoS approaches have competitive or lower levels of stress and crowding than DynNoSlice, thus confirming our research hypothesis, with smaller amounts of movement due to the initial placement. In timesliced graphs, Visone had unsurprisingly the least stress, with the notable exception of InfoVis, where MultiDynNoS and DynNoSlice perform better in terms of both types of stress and crowding. As previously discussed [SAK17], InfoVis is very similar to an event-based data, since there are drastic changes between timeslices as author sets rarely remain stable across consecutive years. On the event-based data, MultiDynNoS and DynNoSlice outperform or match Visone in terms of stress, movement, and crowding. Visone cannot optimize for stress between the timeslices imposed on this naturally expressed event-based data. The video in the supplementary material demonstrates these improvements. The sfdp flat, our baseline, is not able to perform very well in terms of stress on these smaller datasets. However, it is a multilevel algorithm and its strengths are in terms of scalability.

## 5. Conclusion and Future Work

sfdp [Hu05], and is our baseline.

In this paper, we present MultiDynNoS: a multilevel approach for event-based dynamic graph drawing. Our experiment shows an improvement up to an order of magnitude in terms of running time compared to DynNoSlice while retaining its advantages. Future work includes performing a new evaluation on larger datasets that were previously inaccessible to event-based layout techniques.

Table 1: Results of the experiment. |V| and |E| columns report the number of nodes and edges in the flattened graph.  $|E_v|$  reports the number of events in thousands. The **Trend** column visualizes the number of events per timeslice on a scale from 0 to 27% of the total events of the graph. The number of timeslices is reported by the name of the graph in brackets. The **Type** column reports the tested algorithm. The MultiDynNoS variant used is presented as the combination of the initial placement layout (fdp or sfdp) and the coarsening/placement technique used. **T** column reports the algorithm running time in seconds. **Sc.(aling)** column reports the scaling value. Columns **On** and **Off** show the **StressOn** and **StressOff** values. Columns **M** and **C** represent Movement and Crowding respectively; **D** reports the depth of the coarsened hierarchy. MultiDynNoS is implemented in Java 14 and the experiments are run on an i7-8750H CPU with 16GB of RAM.

Timesliced Graphs													
	V	E	$ E_v $	Trend	Type		T (s)	Sc.	On	Off	М	C	D
VanDebunt (7)	39		1 1			sone	0.12	1	1.14	1.46	3.79	0	-
						NoSlice	5.04	0.62	1.23	1.21	3.92	0	-
		32	0.1k			lp flat	0.14	1.61	2.77	2.81	-	0	-
						wi_id	0.48	0.68	1.55	1.62	1.03	0	5
					fdp	is_gr	0.47	0.75	1.03	1.06	0.99	0	3
ek					41	sm_sp	0.46	0.75	1.05	1.08	1.00	0	3
VanI					0.	wi_id	0.56	0.68	1.37	1.39	0.98	0	6
				27%0 0%0	sfdp	is_gr	0.58	0.75	1.09	1.12	0.97	0	3
					Viso	sm_sp	0.58	0.68	1.42	1.48	0.92	0	3
							0.10	1	14.04	14.76	16.36	8	-
(15)	17	93	0.6k	27,		NoSlice	7.58	0.68	16.60	16.57	13.44	1	-
						lp flat	0.15	1.33	26.54	26.52	-	0	-
						wi_id	0.32	0.82	28.40	28.48	2.87	2	6
월				_	fdp	is_gr	0.31	0.82	21.01	20.86	2.95	4	3
Newcomb					J 4	sm_sp	0.32	0.82	22.55	22.39	2.87	1	3
					Q	wi_id	0.42	0.82	27.05	26.94	2.89	2	6
4				%	sfdp	is_gr	0.38	0.82	20.89	20.70	2.82	1	3
					ω	sm_sp	0.39	0.82	21.79	21.71	2.85	2	3
	1,136	2,506	2.8k	27%0 0	Vis	one	77.43	0.46	51.66	52.97	2.14	36	-
				7.7	Dyn	NoSlice	224.93	0.56	30.14	30.19	2.03	2	-
(21)					sfc	lp flat	0.55	1.33	105.29	102.87	-	1,253	-
) <u></u>					0.	wi_id	143.95	0.51	47.26	47.49	0.78	16	7
7.1.8				- % 0	fdp	is_gr	87.79	0.56	28.08	27.79	1.50	4	4
InfoVis						sm_sp	138.95	0.56	28.88	28.65	1.51	4	3
					<u>Q</u>	wi_id	110.00	0.46	51.03	50.97	0.70	36	7
'					gpjs	is_gr	83.00	0.62	28.69	28.59	1.62	2	4
						sm_sp	85.00	0.56	27.21	27.02	1.48	1	3
Event-Based Graphs													
	V	E	$ E_v $	Trend		Type	T(s)	Sc.	On	Off	М	C	D
				27%		sone	0.07	0.68	3.08	2.70	25.46	6	-
				27		NoSlice	2.84	0.51	1.86	1.78	6.64	0	-
0 0					sfc	lp flat	0.18	0.90	2.07	2.02	-	0	-
Rugby (20)	4 -				fdp	wi_id	0.75	0.56	2.18	2.01	1.74	1	5
	12	66	3.1 <i>k</i>		f.	is_gr	1.84	0.56	1.76	1.84	1.25	0	2
						sm_sp	0.52	0.51	2.10	1.94	1.28	0	2
124					sfdp	wi_id	0.88	0.51	2.19	1.97	1.51	1	5
				%	SF	is_gr	1.04	0.513	1.99	1.87	1.11	0	2
						sm_sp	0.77	0.56	2.03	1.95	1.41	0	2
Dialogs (61)	118	501	4.0 <i>k</i>	27%		sone	3.39	0.17	0.62	0.87	5.44	682	-
				[2]		NoSlice	49.53	0.28	0.75	0.90	1.35	0	-
					sic	dp flat	0.21	0.42	0.65	0.69	-	6	1.4
					fdp	wi_id	1.53	0.42	0.53	0.60	0 76	711	14
					ξ	is_gr	5.05	0.35	0.66	0.96	0.76	1	4
						sm_sp	5.49	0.35	0.65	0.91	0.73	0	3
					sfdp	wi_id	1.63	0.42	0.55	0.58	0 71	441	13
				S	sf	is_gr	5.07	0.35	0.64	0.92	0.71	0	4
						sm_sp	5.96	0.31	0.74	0.88	0.64	0	3

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