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## Cyclic polling systems

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CYCLIC POLLING SYSTEMS: LIMITED
SERVICE VERSUS BERNOULLI SCHEDULES
J.P.C. Blanc

FEW 422

CYCLIC POLLING SYSTEMS:
LIMITED SERVICE VERSUS BERNOULLI SCHEDULES

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## Abstract

The power-series algorithm, an iterative numerical technique for the evaluation of the joint queue length distributions for a broad class of multi-queue systems, is applied to cyclic polling systems with a single server. The performances of two service disciplines are compared: limited service and Bernoulli schedules. To a certain extent the properties of the two service disciplines are quite similar, but some striking differences have also been found. The numerical results suggest some general properties of the mean waiting times at the various stations, which may be helpful in deriving better approximations for these quantities.

Keywords: cyclic polling, limited service, Bernoulli schedules, waiting time, power-series algorithm, heavy traffic.

## 1. Introduction

The power-series algorithm is a powerful means for evaluating performance measures of systems consisting of a moderate number of queues. This algorithm is based on power-series expansions of the state probabilities as functions of the load of a system in light traffic. The coefficients of these expansions are computed according to a recursive scheme. The $\varepsilon-a l g o-$ rithm, cf. Wynn [20], is used to improve the convergence of the series in heavy traffic. We refer to Blanc [2,3] for a discussion of the powerseries algorithm. The purpose of this paper is to apply this algorithm to multi-queue models for computer-communication systems, in which a token is passed for access to a single communication channel. In particular, we will compare performance measures of cyclic polling systems with limited service disciplines with those of systems with Bernoulli schedules. On the one hand, limited service seems to be the more preferable discipline since it is a deterministic rule which leads to a bounded number of services during a cycle of the server along the queues. It may be expected that Bernoulli schedules will lead to stronger fluctuations in waiting times. On the other hand, models with Bernoulli schedules seem to be easier to analyse and to optimize, because the parameters of this rule are realvalued, and because this rule requires no memory. Because the power-series algorithm is based on representations of queueing models as Markov-processes, it is indeed easier, in the sense of computing time and storage requirement, to obtain results for systems with Bernoulli schedules than for systems with limited service by means of this algorithm.

There exist only a few analytical results for models with the above mentioned service disciplines. The most general result is the conservation law for systems without switching times, cf. Kleinrock [11], and the pseu-do-conservation law for systems with Bernoulli schedules and with switching times, cf. Tedianto [18]. Cooper [8] derived a set of linear equations, of which the solution determines the mean waiting times in systems with exhaustive service (a limiting case of both disciplines considered in this paper). Models with two queues and 1 -limited service have been solved with the technique which uses Riemann-Hilbert boundary value problems, cf., e.g., Boxma and Groenendijk [6]. Several authors have proposed approximate approaches. Approximations for models with 1-limited service can be found in, e.g., Kühn [12], Arndt and Sulanke [1], Boxma and Meister [4,5] and Srinivasan [15]. Fuhrmann and Wang [9] derive approximations for
the mean waiting times in models with (K-)limited service. Servi [13] discusses an iterative scheme for approximating the mean waiting times for models with Bernoulli schedules. Servi and Yao [14] derive upper and lower bounds for $\mathrm{M} / \mathrm{G} / 1$ vacation models with ( $K$-)limited service and with Bernoulli schedules. Finally, we refer to Takagi [16,17] and Boxma [7] for surveys on polling systems.

The discussion in the present study will be restricted to models with zero switching times, Poisson arrival processes, exponential service times, and infinite buffers. The paper contains several numerical examples and discusses some light and heavy traffic properties of the mean waiting times. The organisation of the paper is as follows. The multi-queue models with cyclic polling will be described in section 2 . Section 3 contains the balance equations for the state probabilities for models with limited service (those for models with Bernoulli schedules can be found in Blanc [3]). Waiting times will be discussed in section 4. Light traffic asymptotes of the mean waiting times have been derived from the recursive scheme of the power-series algorithm; heavy traffic asymptotes have been found with the aid of numerical experiments. Examples of cyclic polling systems for which data are listed in tables will be described in section 5. The final section 6 contains some conclusions based on numerical results. More detailed information on the examples can be found in the appendix.

## 2. The multi-queue model

The system consists of $s$ queues and a single server. Jobs arrive at queue $j$ according to a Poisson process with rate $\lambda_{j}, j=1, \ldots, s$. Each queue may contain an unbounded number of jobs. Service times of jobs at queue $j$ are assumed to be negative exponentially distributed with mean $1 / \mu_{j}, j=1, \ldots, s$. The arrival and service processes are assumed to be independent. The server inspects the queues in a cyclic order ( $1,2, \ldots, s, 1,2$, ...). At each queue jobs are served in order of arrival. The number of jobs which are served during a visit of the server to a certain queue depends on the service discipline at that queue. Performance measures for two service disciplines will be compared in this paper. The first discipline is called limited service. This rule can be described by a vector of positive integers $\bar{K}=\left(K_{1}, \ldots, K_{s}\right)$. When the server visits queue $j$, at most $K_{j}$ jobs are served at this queue. The server proceeds to the nexl queue
when either $K_{j}$ jobs have been served or queue $j$ has become empty, $j=1, \ldots, s$. The second discipline uses Bernoulli schedules. This rule consists of $a$ vector of probabilities $\bar{q}=\left(q_{1}, \ldots, q_{s}\right)$. When the server arrives at a queue, at least one job is served, unless this queue is empty (in this case the server immediately proceeds to the next queue). After the completion of a service at queue $j$ the server starts serving another job at this queue with probability $q_{j}$ if queue $j$ has not yet been emptied; otherwise the server proceeds to the next queue $(j=1, \ldots, s)$. The times which are needed for switching from one queue to another will be neglected in the present study. It should be noted that the power-series algorithm can be used to study a much broader class of polling systems, e.g. with Coxian switching times, with Coxian service times, with finite buffers and with polling according to a table or with random polling, see also Blanc [3].

It holds for both service disciplines that the distribution of the total amount of work in the system is equal to that of an M/G/1 system with arrival rate $\wedge$ and with $\beta_{1}$ and $\beta_{2}$ as first two moments of the service time distribution:

$$
\begin{equation*}
\wedge:=\sum_{j=1}^{S} \lambda_{j}, \quad \beta_{1}:=\sum_{j=1}^{S} \frac{\lambda_{j}}{\wedge} \frac{1}{\mu_{j}}, \quad \beta_{2}:=2 \sum_{j=1}^{S} \frac{\lambda_{j}}{\wedge}\left(\frac{1}{\mu_{j}}\right)^{2} . \tag{2.1}
\end{equation*}
$$

Therefore, the traffic intensities $\rho$ of the polling systems are defined as

$$
\begin{equation*}
\rho:=\wedge \beta_{1}=\wedge \sum_{j=1}^{s} \frac{\lambda_{j}}{\wedge} \frac{1}{\mu_{j}}=\sum_{j=1}^{s} \frac{\lambda_{j}}{\mu_{j}}, \tag{2.2}
\end{equation*}
$$

and a necessary and sufficient condition for ergodicity of the systems is $p<1$. It will be assumed throughout that the systems are in steady state. Finally, we introduce the following load-independent quantities:

$$
\begin{equation*}
a_{j}:=\lambda_{j} / \rho, \quad \eta_{j}:=a_{j} / \mu_{j}, \quad j=1, \ldots, s . \tag{2.3}
\end{equation*}
$$

Note that definitions (2.1), (2.2) and (2.3) imply the following relations

$$
\begin{equation*}
\sum_{j=1}^{s} \eta_{j}=1, \quad \sum_{j=1}^{s} a_{j}=\frac{\Lambda}{\rho}=\frac{1}{\beta_{1}}, \quad \beta_{2}=2 \beta_{1} \sum_{j=1}^{s} \frac{\eta_{j}}{\mu_{j}} . \tag{2.4}
\end{equation*}
$$

## 3. Balance equations for limited service models

The power-series algorithm has been described in Blanc [3] for systems with Bernoulli schedules. Below we will only present the balance equations for the state probabilities of systems with limited service. The recursive scheme of the power-series algorithm can be derived from these equations in the same way as in Blanc [3]. Let $N_{j}$ denote the number of jobs in queue $j$ (waiting or being served), $j=1, \ldots, s$. In order to transform the queue length process of limited service systems into a Markov process we introduce a polling table and a supplementary variable $H$, indicating the actual position in the table. The polling table is described as follows. Let $L:=\sum_{j=1}^{S} K_{j}$ be the length of the table. The mapping $\ell(h)$ from table entry to queue number is defined by

$$
\begin{equation*}
\ell(h)=j, \quad \text { if } \sum_{i=1}^{j-1} K_{i}<h \leq \sum_{i=1}^{j} K_{i}, \text { for } j=1, \ldots, s, \quad h=1, \ldots, L, \tag{3.1}
\end{equation*}
$$

and it is continued as a periodic function by the convention

$$
\begin{equation*}
\ell(h+k L)=\ell(h), \quad h \in\{1, \ldots, L\}, \quad k \in Z \tag{3.2}
\end{equation*}
$$

The value of the variable $H$ is increased by one whenever a service has been completed or when queue $\ell(H)$ is empty, unless the whole system has become empty; in the latter case the value of $H$ is set and kept equal to 1 until a new arrival occurs. The value of $\ell(H)$ determines the queue to which the server is attending. Let $\bar{n}=\left(n_{1}, \ldots, n_{s}\right)$ be a vector with nonnegative integer entries. The state probabilities are defined as follows:

$$
\begin{equation*}
p(\bar{n}, h):=\operatorname{Pr}\left\{N_{j}=n_{j}, j=1, \ldots, s ; H=h\right\}, \quad \bar{n} \in \mathbb{N}^{S}, \quad h=1, \ldots, L \tag{3.3}
\end{equation*}
$$

Let $I\{E\}$ stand for the indicator function of the event $E$, and let $\bar{e}_{j}$ be a vector with zero entries except an entry of one at the $j^{\text {th }}$ position ( $j=1, \ldots, s$ ). The balance equations for the state probabilities (3.3) in models with a limited service discipline are readily verified to be, for $h=1, \ldots, L, \bar{n} \in \mathbb{N}^{s}$,

$$
\begin{align*}
& {\left[\rho \sum_{j=1}^{s} a_{j}+\mu_{\ell(h)}\right] p(\bar{n}, h)=p \sum_{j=1}^{s} a_{j} p\left(\bar{n}-\bar{e}_{j}, h\right) I\left\{n_{j}>0 ; n_{j}>1 \text { if } j=\ell(h)\right\}} \\
& \quad+\sum_{i=1}^{L} \mu_{\ell(h-i)} p\left(\bar{n}+\bar{e}_{\ell(h-i)} ; h-i\right) I\left\{n_{\ell(\nu)}=0, \nu=h-i+1, \ldots, h-1\right\} \\
& \quad+a_{\ell(h)} p p(\overline{0}, 1) I\left\{\bar{n}=\bar{e}_{\ell(h)} \wedge \ell(i) \neq \ell(h), i=1, \ldots, h-1\right\}, \quad{ }^{n} \ell(h)>0 ;  \tag{3.4}\\
& \rho \sum_{j=1}^{s} a_{j} p(\overline{0}, 1)=\sum_{h=1}^{L} \mu_{\ell(h)} p\left(\bar{e}_{\ell(h)}, h\right) .
\end{align*}
$$

Note that a state $(\bar{n}, \mathrm{~h})$ with $n_{\ell(h)}=1$ can be entered through an arrival at queue $h$ only if $\bar{n}=\bar{e}_{\ell(h)}$ and if $h$ is the first entry on the polling table with the value $\ell(h), h=1, \ldots, L$. Further, it should be noted that the balance equations (3.4) and (3.5) are valid for models with arbitrary polling tables, i.e. for arbitrary surjective mappings $\ell:\{1, \ldots, L\} \rightarrow\{1, \ldots, s\}$. The examples in this paper, however, are restricted to models with limited service, cf. (3.1). The reader is referred to Blanc [3] for details concerning the derivation of a recursive scheme from the balance equations, the computation of moments of the joint queue length distribution, and the application of the $\varepsilon$-algorithm.

## 4. Waiting times

This section is concerned with a discussion of the stationary distributions of the waiting times $W_{j}$ (excluding service times) of jobs arriving at queue $j(j=1, \ldots, s)$, for both limited service and Bernoulli schedules. Firstly, some general relations will be reviewed. Then, two terms of the power-series expansions of the mean waiting times in light traffic will be derived from the recursive scheme of the power-series algorithm. Finally, several properties concerning the heavy traffic behaviour of the mean waiting times, which have been found on the basis of numerical data, will be discussed.

The power-series algorithm computes the joint queue-length distribution. The waiting time distributions are related to this distribution through (Blanc [3]):

$$
\begin{equation*}
E\left\{z^{N_{j}}\right\}=\frac{1}{1+(1-z) n_{j} \rho} E\left\{e^{-a_{j} \rho(1-z) W_{j}}\right\}, \quad|z| \leq 1, \quad j=1, \ldots, s \tag{4.1}
\end{equation*}
$$

Let $W$ be the waiting time of an arbitrary job. Then, with (2.2) and (2.3),

$$
\begin{equation*}
E\{W\}=\sum_{j=1}^{s} \frac{\lambda_{j}}{\wedge} E\left\{W_{j}\right\}=\beta_{1} \sum_{j=1}^{s} a_{j} E\left\{W_{j}\right\} \tag{4.2}
\end{equation*}
$$

We recall that the mean waiting times at the various queues of a polling system satisfy the following conservation law, cf. Kleinrock [11], Boxma [7].

$$
\begin{equation*}
\sum_{j=1}^{S} \eta_{j} E\left\{W_{j}\right\}=\frac{\rho}{1-\rho} \sum_{j=1}^{S} \frac{\eta_{j}}{\mu_{j}}=\frac{\rho}{1-\rho} \frac{\beta_{2}}{2 \beta_{1}} \tag{4.3}
\end{equation*}
$$

Next, we discuss the light traffic behaviour of the mean waiting times. For this purpose we introduce for $j=1, \ldots, s, \bar{q}=\left(q_{1}, \ldots, q_{s}\right)$, the quantities:

$$
\equiv(j):=2 \sum_{i=1}^{s-1} \eta_{j+i} \sum_{\nu=0}^{i-1} \frac{\eta_{j+\nu}}{\mu_{j+\nu}}, \quad \Psi(j, \bar{q}):=2 \sum_{i=1}^{s} q_{i} \frac{\eta_{i}^{2}}{\mu_{i}}-2 q_{j} \frac{\eta_{j}}{\mu_{j}} ;
$$

here, indices exceeding s should be read modulo s.

Theorem 1a. For cyclic polling systems with Bernoullt schedules $\bar{q}$ and with exponential service time distributions it holds for $j=1, \ldots, s$, that,

$$
E\left\{W_{j}\right\}=\rho \frac{\beta_{2}}{2 \beta_{1}}+p^{2}\left[\eta_{j} \frac{\beta_{2}}{2 \beta_{1}}+\Xi(j)+\Psi(j, \bar{q})\right]+O\left(\rho^{3}\right), \quad \text { as } p \downarrow 0
$$

Theorem 1b. For cyclic polling systems with limited service discipline $\bar{K}$ and with exponential service time distributions the power-series expansions (4.4) hold with $q_{j}$ replaced by $I\left\{K_{j} \geq 2\right\}, j=1, \ldots, s$.

Proof. In order to derive these light traffic asymptotes we determine first the coefficients of the power-series expansions of the state probabilities and the mean queue lengths up to the $3^{\text {rd }}$ power of $\rho$ according to the recurrence relations of the power-series algorithm (see Blanc [3]). This leads for Bernoulli schedules to: for $j=1, \ldots, s$, as $\rho \downarrow 0$,

$$
E\left\{N_{j}\right\}=\eta_{j} p+a_{j} \rho^{2} \frac{\beta_{2}}{2 \beta_{1}}+a_{j} p^{3}\left[\eta_{j} \frac{\beta_{2}}{2 \beta_{1}}+\equiv(j)+\Psi(j, \bar{q})\right]+O\left(\rho^{4}\right)
$$

Then, the coefficients of the power-series expansions of the mean waiting times follow with the aid of Little's formula, cf. (4.1).

The first term of the coefficient of $p^{2}$ in the power-series expansions of the mean waiting times is independent of the order in which the queues are placed and of the service disciplines at the queues. The term $\Xi$ (.) reflects the influence of the order in which the server visits the queues, and the term $\Psi(.,$.$) depends on the service disciplines at the queues. Note$ that the coefficients of the power-series expansions of the mean waiting times up to the $m^{\text {th }}$ power of $\rho, m=1,2, \ldots$, are the same for all $K_{i} \geq m$, $i=1, \ldots, s$, in systems with limited service. It is possible, but increasingly tedious, to determine more coefficients of the power-series expansions of the mean waiting times in a similar way. The appendix contains a discussion of the above mentioned functions and properties for some special cases.

For the description of heavy traffic properties we introduce the limits:

$$
\begin{align*}
& \omega_{j}:=\lim _{\rho \uparrow 1}^{\lim }(1-\rho) E\left\{W_{j}\right\}, \quad j=1, \ldots, s, \quad \omega_{0}:=\lim _{\rho \uparrow 1}^{\lim }(1-\rho) E\{W\},  \tag{4.5}\\
& x_{j}:=\lim _{\rho \uparrow 1}^{\lim }\left[E\left\{W_{j}\right\}-\frac{\omega_{j}}{(1-\rho)}\right], \quad j=1, \ldots, s, \quad x_{0}:=\lim _{\rho \uparrow 1}^{\lim }\left[E\{W\}-\frac{\omega_{0}}{(1-\rho)}\right] . \tag{4.6}
\end{align*}
$$

The conservation law (4.3) and relation (4.2) imply that

$$
\begin{array}{ll}
\sum_{j=1}^{s} \eta_{j} \omega_{j}=\frac{\beta_{2}}{2 \beta_{1}}, & \omega_{0}=\beta_{1} \sum_{j=1}^{s} a_{j} \omega_{j}, \\
\sum_{j=1}^{S} \eta_{j} x_{j}=-\frac{\beta_{2}}{2 \beta_{1}}, & x_{0}=\beta_{1} \sum_{j=1}^{s} a_{j} x_{j} .
\end{array}
$$

Theorem 2. If the service discipline is exhaustive at each queue, i.e. $q_{j}=1, j=1, \ldots, s$, then the limits defined in (4.5) are given by:

$$
\begin{equation*}
\omega_{j}=\frac{1-\eta_{j}}{\sum_{i=1}^{s} \eta_{i}\left(1-\eta_{i}\right)} \frac{\beta_{2}}{2 \beta_{1}}, \quad j=1, \ldots, s \tag{4.8}
\end{equation*}
$$

Proof. From the results of Cooper [8] it can be deduced that when the service discipline is exhaustive at each queue, it holds that $\omega_{j}=\delta_{j-1}(1)$, $j=1, \ldots, s$, where the quantities $\delta_{j}(k), j=1, \ldots, s, k=1, \ldots, s-1$, satisfy the following set of $s \times(s-1)$ linear equations (read $\delta_{j+s}(k)$ for $\delta_{j}(k)$ and $\eta_{j+s}$ for $\eta_{j}$ whenever $j<1$ ):

$$
\begin{aligned}
\left(1-\eta_{j+1}\right) \delta_{j}(k)= & \sum_{h=0}^{s-1-k} \frac{\eta_{j-h}^{2}}{1-\eta_{j-h}} \delta_{j-1-h}(1)+\sum_{h=0}^{s-2-k} \eta_{j-h} \delta_{j-1-h}(k+1+h) \\
& +\sum_{h=0}^{s-1-k} \eta_{j-h} \delta_{j-1-h}(2+h) I\{h<s-2\} .
\end{aligned}
$$

It is rather tedious, but straightforward, to verify with the aid of (2.4) that solutions of this set of equations are of the form

$$
\delta_{j}(k)=c \sum_{i=0}^{s-1-k} \eta_{j-i}, \quad j=1, \ldots, s, \quad k=1, \ldots, s-1
$$

This implies that the quotient $\omega_{j} /\left(1-\eta_{j}\right)$ has the same value for each $j$, $j=1, \ldots . s$. The constant $C$ can be determined with the aid of the conservation law (4.3). a

The proof and the assertion of theorem 2 remain valid for general service time distributions if $\beta_{2}$ is read as the second moment of the job-averaged service time distribution, cf. (2.1), because the set of equations for $\delta_{j}(k), j=1, \ldots, s, k=1, \ldots, s-1$, only depends on the mean service times, and because the conservation law (4.3) holds for general service time distributions.

Property 1. The limit $\omega_{j}, j=1, \ldots, s$, is positive if and only if (in the cases of limited service and Bernoulli schedules, respectively)

$$
\begin{equation*}
a_{j} / K_{j}=\max _{i=1, \ldots, s}\left\{a_{i} / K_{i}\right\}, \quad a_{j}\left(1-q_{j}\right)=\max _{i=1, \ldots, s}\left\{a_{i}\left(1-q_{i}\right)\right\} \tag{4.9}
\end{equation*}
$$

Remark. When the lefthand sides of (4.9) are close to, but not equal to, the maximum at the righthand sides of these relations for some $j$, $j=1, \ldots$ s, the limit $\omega_{j}$ is zero, but $E\left\{W_{j}\right\}$ has a large finite limit as $p \uparrow$ 1. That only the arrival rates, and not the service rates, play a role in property 1, can be explained by the fact that a certain (integer) number of jobs is served during each cycle of the server along the queues according to the limited service discipline as well as to the Bernoulli schedules. If service disciplines would be considered in which the server spends a certain amount of time at each queue during a cycle and in which service of a job can be interrupted and resumed in a later cycle, then it might be expected that a similar property as property 1 holds, but with the arrival rates replaced by the relative loads at the queues. Note that the relations (4.9) hold for each queue in systems with exhaustive service at each queue $\left(K_{j}=\infty\right.$ or $\left.q_{j}=1, j=1, \ldots, s\right)$. o

As a consequence of property 1 the $\varepsilon$-algorithm which is being used to accelerate the convergence of the power-series occuring in the algorithm, should not be modified as described in Blanc [3] for moments of the marginal queue length distributions at queues where (4.9) does not hold. Another implication is that approximations for mean waiting times in polling systems, which do not possess property 1 will behave poorly under heavy traffic circumstances. For instance, the approximations for the mean waiting times in Boxma and Meister [4] have the right heavy traffic limits in case of exhaustive service (formula (17) in [4]) according to theorem 2. but they do not have the proper heavy traffic behaviour in case of 1limited service (formula (20) in [4]) according to property 1.

Corollary 1. If there exists a queue $j$ such that for each $i, i=1, \ldots, s$, $i \neq j$,

$$
\begin{equation*}
a_{i} / K_{i}<a_{j} / K_{j}, \quad \text { respectively } \quad a_{i}\left(1-q_{i}\right)<a_{j}\left(1-q_{j}\right) \text {, } \tag{4.10}
\end{equation*}
$$

then property 1 implies that $\omega_{i}=0$ for each $i, i=1, \ldots, s, i \neq j$, and that

$$
\begin{equation*}
\omega_{j}=\frac{1}{\eta_{j}} \frac{\beta_{2}}{2 \beta_{1}}, \quad \omega_{0}=\frac{1}{2} \mu_{j} \beta_{2} \tag{4.11}
\end{equation*}
$$

Remark. The above corollary is a consequence of property 1 and of (4.7). Watson [19] derived a similar result for systems with 1-1imited service and non-negligible switching times. व

Property 2. If the system consists of two queues, if the service discipline is Bernoulli at each queue, and if $a_{1}\left(1-q_{1}\right)>a_{2}\left(1-q_{2}\right)$, then:

$$
\begin{equation*}
\omega_{2}=0, \quad x_{2}=\frac{\mu_{1} \mu_{2} \frac{\beta_{2}}{2 \beta_{1}}+\eta_{2}\left[\mu_{2}\left(1-q_{2}\right)-\mu_{1} q_{1}\right]}{a_{1}\left(1-q_{1}\right)-a_{2}\left(1-q_{2}\right)} \frac{1}{\mu_{2}} \tag{4.12}
\end{equation*}
$$

and the quantities $\omega_{1}, \omega_{0}, x_{1}$, and $x_{0}$ follow from (4.7).

Remark. Property 2 has been found on the basis of numerical experiments with several values of the parameters $a_{1}, a_{2}, \mu_{1}, \mu_{2}, q_{1}$ and $q_{2}$. Further, this property agrees with the following observations. Corollary 1 implies that $\omega_{2}=0$. The denominator of $x_{2}$ vanishes when $a_{1}\left(1-q_{1}\right)=a_{2}\left(1-q_{2}\right)$, which is in agreement with property 1 . In the case $q_{1}=0, q_{2}=1$, jobs at queue 2 have non-preemptive priority over jobs at queue 1. For this model it is known (see Jaiswal [10], §v.3) that for $0<p<1$,

$$
E\left\{W_{2}\right\}=(1-\rho) E\left\{W_{1}\right\}=\frac{\rho}{1-n_{2} \rho} \frac{\beta_{2}}{2 \beta_{1}},
$$

from which it is readily seen that property 2 holds in this case. a

We did not find a general result similar to property 2 for systems with two queues and limited service, except that $\omega_{2}$ vanishes if $a_{1} / K_{1}>a_{2} / K_{2}$. by corollary 1. Only in the case that $K_{2}=\infty$ we have found that

$$
\begin{equation*}
x_{2}=\frac{1}{\eta_{1}} \frac{\beta_{2}}{2 \beta_{1}}+\frac{1}{2 \mu_{1}}\left(K_{1}-1\right) . \tag{4.13}
\end{equation*}
$$

Note that in this case $\left(K_{2}=\infty\right.$ or $\left.q_{2}=1\right)$ the limits $x_{2}$ in (4.12) and (4.13) agree when $K_{1}=\left(1+q_{1}\right) /\left(1-q_{1}\right)$.

## 5. Numerical examples

This section contains descriptions of cyclic polling systems for which data have been generated with the aid of the power-series algorithm together with the $\varepsilon$-algorithm (cf. Wynn [20], Blanc [3]).

In table 1 the standard deviation $\sigma\{W\}$ of the waiting times has been listed for symmetrical systems with limited service ( $\left.K_{j}=K, j=1, \ldots, s\right)$. Table 2 contains the standard deviation of the waiting times for symmetrical systems with Bernoulli schedules ( $q_{j}=q, j=1, \ldots, s$ ). The mean waiting time follows directly from (4.3) for symmetrical systems, and does not depend on the service discipline. In both tables, $\mu_{j}=1$ for $j=1, \ldots$ s.

Tables $3 a$ and $4 a$ show in which way the waiting time distributions depend on the arrival rates and the service rates, respectively. Both tables are concerned with models with three queues, with traffic intensity $\rho=0.90$ and with relative loads in the proportion of $\eta_{1}: \eta_{2}: \eta_{3}=1: 2: 3$. In table 3a, the service rates at the three queues are equal ( $\mu_{j}=1, j=1,2,3$ ), and the arrival rates are $\lambda_{1}=0.15, \lambda_{2}=0.30$, and $\lambda_{3}=0.45$. In table $4 a$, the arrival rates at the three queues are equal $\left(\lambda_{j}=0.30, j=1,2,3\right)$, and the service rates are $\mu_{1}=2, \mu_{2}=1, \mu_{3}=2 / 3$. Because the mean waiting time $\mathrm{E}\{\mathrm{W}\}$ is equal to 9.000 for the models in table 3 a , independently of the service discipline, cf. (4.2), (4.3), this quantity has been omitted. The value of the righthand side of the conservation law (4.3) is equal to 10.500 for the models in table 4 a . In both tables the data on the lefthand side concern systems with a limited service discipline ( $\mathrm{K}_{1}, \mathrm{~K}_{2}, \mathrm{~K}_{3}$ ) and those on the righthand side concern systems with Bernoulli schedules $\left(q_{1}, q_{2}, q_{3}\right)$. In these and the following tables, "dcp" stands for service discipline. The server visits the queues in the usual order ( $1,2,3,1, \ldots$ ) in the models in the tables $3 a$ and $4 a$. Tables $3 b$ and $4 b$ contain data for the same models as tables $3 a$ and $4 a$ respectively, but the order in which the server inspects the queues has been reversed $(1,3,2,1, \ldots)$. The routing of the server does not seem to have a major influence on the performance measures, compared with the influence of the parameters of the service disciplines. Still, altering the polling order may lead to differences in mean waiting times of more than $10 \%$.

Tables $5 a$ and $5 b$ show the influence of a relatively heavily loaded queue on the mean waiting times at queues which are four times less heavily loaded. The parameters of the system are in the case of $s$ queues, $s=2, \ldots, 6: \quad \mu_{1}=1, \mu_{j}=2, j=2, \ldots, s ; a_{1}=2 a_{j}=\frac{4}{s+3}, j=2, \ldots, s$ (hence,
$\left.\beta_{2} / \beta_{1}=\frac{s+7}{s+3}\right)$; and $\rho=0.75$. Table 6 contains data for the same models as table 5a, but for a traffic intensity of $p=0.95$. The examples concern only service disciplines with $K_{j}=K_{2}, j=2, \ldots, s$, or $q_{j}=q_{2}, j=2, \ldots, s$; therefore, the service discipline is indicated either by ( $K_{1}, K_{2}$ ) or by $\left(q_{1}, q_{2}\right)$. Note that the differences in mean waiting times of the lightly loaded queues are not negligible, although their arrival and service rates and their service disciplines are the same. In the cases of 4,5 and 6 queues the values of $E\left\{W_{3}\right\}, \ldots, E\left\{W_{S-1}\right\}$ lie, in this order, in between those of $E\left\{W_{2}\right\}$ and $E\left\{W_{S}\right\}$ for all considered service disciplines.

Although we do not have the disposal of bounds on errors for data generated by the power-series algorithm together with the $\varepsilon$-algorithm, we estimate on the basis of differences between successive terms that relative errors are below $1 \%$ for almost all quantities listed in the tables, and even below $0.01 \%$ for most quantities. In general, errors increase with increasing traffic intensity, with increasing number of queues, with increasing length of the polling table (L), with increasing differences in the arrival rates and the service rates, and when equality in (4.9) is approximated (i.e. when the mean waiting time at one or more queues possesses a large but finite heavy traffic limit). Also, errors for standard deviations are usually larger than those for averages.

## 6. Conclusions

When comparing data for the waiting time distributions in systems with a limited service discipline with those with Bernoulli schedules we arrive at the following conclusions. Tables 1 and 2 show that Bernoulli schedules lead to higher variances of the waiting times than limited service disciplines in symmetrical systems. In fact, the standard deviation of the waiting times seems to be a convex function of $K$ for models with limited service disciplines and a concave function of $q$ for models with Bernoulli schedules. It is interesting to note that while the standard deviation of the waiting times in symmetrical systems with exhaustive service $(q=1)$ is larger than that in symmetrical systems with 1-limited service $(q=0)$ when the number of queues is small ( $s \leq 4$ ), this property does no longer hold when the number of queues increases. The limited service disciplines and the Bernoulli schedules agree in that the mean waiting times pass globally through similar trajectories (though continuously in case of Bernoulli schedules and with jumps in case of limited service),
when we consider them as functions of one parameter, e.g., $K_{j}=K, q_{j}=q$ $(j=1, \ldots, s)$ or as function of $K_{j}$ and $q_{j}$ for some $j, j=1, \ldots, s$, while the other queues have exhaustive or 1-limited disciplines. But there seems to be no general relationship between the parameters of the Bernoulli and of the limited service schedules, which give approximately the same position on these trajectories. See tables 3, 4,5,6, and compare (4.10), which suggests a relation between $K_{j}$ and $1 /\left(1-q_{j}\right)$, with the observation below (4.13). It should be noted that mean waiting times are not in every case monotonous functions of the parameters of the service discipline; see for instance table $3 a$ where $E\left\{W_{1}\right\}$ is larger for $q_{1}=0.90$ then for both $q_{1}=0.00$ and $q_{1}=1.00$ when $q_{j}$ is related to $q_{1}$ by $a_{j}\left(1-q_{j}\right)=a_{1}\left(1-q_{1}\right)$, $j=2,3$. A general property suggested by the examples, and supported by theorems 1 a and 1 b , is that $\mathrm{E}\left\{\mathrm{W}_{\mathrm{j}}\right\}$ is minimal over all disciplines $\overline{\mathrm{K}}$ and $\overline{\mathrm{q}}$, for fixed arrival and service rates, when $q_{j}=0 \quad\left(K_{j}=1\right)$ and $q_{i}=1 \quad\left(K_{i}=\infty\right)$, $i \neq j, i=1, \ldots, s$, and maximal in the reversed case $q_{j}=1$ and $q_{i}=0, i \neq j$, $i=1, \ldots, s(j=1, \ldots, s)$. Finally, we note that examples show that the ordering of the mean waiting times in certain models is not the same for all values of $p$ (compare, e.g., the ordering for $q_{1}=q_{2}=0.90$ in the tables 5 a and 6 ; see also table A.3). This is supported by theorem 1a (1b) and property 1.

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Table 1. Standard deviation of the waiting time for symmetrical systems with limited service disciplines.

| s | K | $\rho=.5$ | $\rho=.75$ | $\rho=.9$ | $\rho=.95$ | $\rho=.98$ | s K | $\rho=.5$ | $\rho=.75$ | $\rho=.9$ | $\rho=.95$ | $\rho=.98$ |  |
| :--- | ---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 2 | 1 | 1.8350 | 4.246 | 11.187 | 22.67 | 57.09 | 3 | 1 | 1.8805 | 4.434 | 11.87 | 24.2 | 61.2 |
| 2 | 2 | 1.8301 | 4.237 | 11.174 | 22.66 | 57.08 | 3 | 2 | 1.8739 | 4.410 | 11.82 | 24.1 | 61.1 |
| 2 | 3 | 1.8320 | 4.231 | 11.159 | 22.64 | 57.05 | 3 | 3 | 1.8788 | 4.396 | 11.77 | 24.1 | 61.0 |
| 2 | 4 | 1.8388 | 4.230 | 11.145 | 22.62 | 57.03 | 3 | 4 | 1.8888 | 4.394 | 11.73 | 24.0 | 60.9 |
| 2 | 5 | 1.8479 | 4.234 | 11.133 | 22.60 | 57.00 | 3 | 5 | 1.8991 | 4.401 | 11.69 | 23.9 | 60.8 |
| 2 | 6 | 1.8575 | 4.243 | 11.124 | 22.58 | 56.97 | 3 | 6 | 1.9076 | 4.415 | 11.67 | 23.9 | 60.7 |
| 2 | 7 | 1.8664 | 4.254 | 11.118 | 22.56 | 56.95 | 3 | 7 | 1.9138 | 4.432 | 11.64 | 23.8 | 60.4 |
| 2 | 8 | 1.8741 | 4.269 | 11.114 | 22.54 | 56.92 | 3 | 8 | 1.9180 | 4.452 | 11.63 | 23.7 | 60.3 |
| 2 | 9 | 1.8805 | 4.285 | 11.113 | 22.53 | 56.89 | 3 | $\infty$ | 1.9254 | 4.610 | 12.41 | 25.3 | 64.1 |
| 2 | 10 | 1.8857 | 4.302 | 11.115 | 22.51 | 56.85 |  |  |  |  |  |  |  |
| 2 | 12 | 1.8927 | 4.337 | 11.126 | 22.49 | 56.78 | 4 | 1 | 1.9057 | 4.549 | 12.31 | 25.2 | 64. |
| 2 | 14 | 1.8968 | 4.371 | 11.146 | 22.47 | 56.74 | 4 | 2 | 1.8973 | 4.507 | 12.22 | 25.1 | 64. |
| 2 | 16 | 1.8990 | 4.403 | 11.173 | 22.46 | 56.68 | 4 | 3 | 1.9033 | 4.483 | 12.11 | 24.9 | 63. |
| 2 | 18 | 1.9002 | 4.432 | 11.207 | 22.46 | 56.62 | 4 | 4 | 1.9131 | 4.478 | 12.00 | 24.6 | 63. |
| 2 | 20 | 1.9008 | 4.456 | 11.245 | 22.46 | 56.59 | 4 | 5 | 1.9214 | 4.485 | 11.90 | 24.3 | 62. |
| 2 | $\infty$ | 1.9014 | 4.569 | 12.362 | 25.29 | 64.03 | 4 | $\infty$ | 1.9345 | 4.624 | 12.42 | 25.3 | 64. |

Table 2. Standard deviation of the waiting time for symmetrical systems
with Bernoulli schedules.

| q | s | $p=.5$ | $\rho=.75$ | $p=.9$ | $p=.95$ | $p=.98$ | s | $p=.5$ | $p=.75$ | $p=.9$ | $p=.95$ | 98 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| . 00 | 2 | 1.8350 | 4.246 | 11.187 | 22.67 | 57.09 | 5 | 1.9213 | 4.624 | 12.62 | 25.9 | 65.9 |
| . 25 | 2 | 1.8448 | 4.289 | 11.340 | 23.01 | 57.99 | 5 | 1.9283 | 4.665 | 12.80 | 26.4 |  |
| . 50 | 2 | 1.8575 | 4.347 | 11.547 | 23.47 | 59.22 | 5 | 1.9352 | 4.706 | 13.01 | 26.9 | 68.8 |
| . 75 | 2 | 1.8748 | 4.430 | 11.849 | 24.15 | 61.01 | 5 | 1.9407 | 4.734 | 13.22 | 27.6 | 70.9 |
|  | 2 |  | 4.502 | 12.116 | 24.74 | 62.59 | 5 | 1.9415 | 4.716 | 13.19 | 27.7 | 72.0 |
| . 98 | 2 | 1.8987 | 4.554 | 12.230 12.307 | 25.00 | 63.26 | 5 | 1.9409 | 4.689 | 13.02 | 27.4 | 71.5 |
| . 99 | 2 | 1.9000 | 4.562 | 12.334 | 25.23 | 63.86 | 5 | 1.9401 1.9398 | 4. | 12.78 12.63 | 26.7 26.2 | 70.0 68.4 |
| 1.0 | 2 | 1.9014 | 4.569 | 12.362 | 25.29 | 64.03 | 5 | 1.9394 | 4.630 | 12.43 | 25.4 | 64.1 |
| .00 | 3 | 1.8805 | 4.434 | 11.865 | 24.20 | 61.17 | 6 | 1.9316 | 4.677 | 12.84 | 26.5 | 67.4 |
| . 25 | 3 | 1.8904 | 4.483 | 12.055 | 24.63 | 62.35 | 6 | 1.9373 | 4.712 | 13.02 | 26.9 | 68.7 |
| . 50 | 3 | 1.9019 | 4.542 | 12.296 | 25.20 | 63.91 | 6 | 1.9425 | 4.744 | 13.20 | 27.4 | 70.2 |
| . 75 | 3 | 1.9149 | 4.612 | 12.604 | 25.95 | 66.05 | 6 | 1.9457 | 4.756 | 13.34 | 27.9 | 72.0 |
| . 90 | 3 | 1.9224 | 4.646 | 12.782 | 26.46 | 67.70 | 6 | 1.9452 | 4.724 | 13.24 | 27.9 | 72.5 |
| . 9 | 3 | 1.9243 | 4.644 | 12.773 | 26.51 | 68.12 | 6 | 1.9441 | 4.692 | 13.03 | 27.4 | 71.8 |
|  | 3 | 1.9251 | 4.630 | 12.665 | 26.28 | 67.83 | 6 | 1.9431 | 4.662 | 12.77 | 26.7 | 70.1 |
| 199 | 3 | 1.9252 1.9254 | 4.622 4.610 | 12.574 12.408 | 26.02 | 67.16 | 6 | 1.9428 | 4.649 | 12.63 | 26.2 | 68.4 |
| 1.0 | 3 | 1.9254 | 4.610 | 12.408 | 25.33 | 64.07 | 6 | 1.9423 | 4.634 | 12.43 | 25.4 | 64.1 |
| . 00 | 4 | 1.9057 | 4.549 | 12.307 | 25.21 | 63.93 | 7 | 1.9388 | 4.713 | 13.0 | 26.8 | 68. |
| . 25 | 4 | 1.9142 | 4.594 | 12.501 | 25.68 | 65.22 | 7 | 1.9434 | 4.742 | 13.1 | 27.2 | 69. |
| . 50 | 4 | 1.9233 | 4.645 | 12.734 | 26.25 | 66.86 | 7 | 1.9472 | 4.763 | 13.3 | 27.6 | 71. |
| . 75 | 4 | 1.9319 | 4.693 | 12.993 | 26.95 | 68.99 | 7 | 1.9489 | 4.765 | 13.4 | 27.9 | 72. |
| . 90 | 4 | 1.9352 | 4.696 | 13.059 | 27.28 | 70.39 | 7 | 1.9474 | 4.726 | 13.2 | 27.8 | 72. |
| . 95 | 4 | 1.9353 | 4.677 | 12.958 | 27.13 | 70.43 | 7 | 1.9462 | 4.693 | 13.0 | 27.4 | 71. |
|  | 4 | 1.935 | 4.652 | 12.757 | 26.63 | 69.46 | 7 | 1.9452 | 4.663 | 12.8 | 26.6 | 67. |
| 1.0 | 4 | 1.9345 | 4.624 | 12.422 | 25.35 | 68.25 64.09 | 7 | 1.9448 1.9444 | 4.650 4.637 | 12.6 | 26.2 | 65. |

Table 3a. The waiting time distributions for systems with three queues, unequal arrival rates and equal service rates.

| dcp | $\mathrm{E}\left\{\mathrm{W}_{1}\right.$ | $\mathrm{w}_{2}$ | W3 | $\sigma\{\mathrm{W}\}$ |  | dcp |  | \} | $\mathrm{E}\left\{\mathrm{W}_{2}\right\}$ | E\{ | $\sigma\{\mathrm{W}\}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 111 | 2.445 | 5.048 | 13.820 | 13.7 | . 25 | 25 | 25 | 2.856 | . 4 |  |  |
| 222 | 2.790 | 5.141 | 13.643 | 13.646 | . 50 | . 50 | 50 | 3.570 | 6.066 | 12.766 | 13.261 |
| 333 | 3.216 | 5.264 | 13.419 | 13.534 | . 75 | . 75 | 75 | 5.146 | 7.306 | 11.414 | 12.794 |
| 444 | 3.671 | 5.406 | 13.173 | 13.406 | . 90 | . 90 | . 90 | 7.578 | 8.733 | 9.652 | 12.526 |
| 666 | 4.583 | 5.727 | 12.666 | 13.253 | . 95 | . 95 | 95 | 9.186 | 9.361 | 8.697 | 12.575 |
| 888 | 5.439 | 6.077 | 12.136 | 12.926 | . 99 | . 99 | 99 | 11.148 | 9.791 | 7.756 | 12.732 |
| 12 | 11.828 | 9.158 | 7.952 | 12.005 | . 00 | . 50 | 67 | 10.936 | 9.476 | 8.037 | 12.326 |
|  | 11.892 | 9.251 | 7.868 | 11.938 | . 50 | 75 | 83 | 11.457 | 9.597 | 7.783 | 12.740 |
| $\begin{aligned} & 369 \\ & 4812 \end{aligned}$ | 11.939 | 9.345 | 7.790 | 11.899 | . 80 | 90 | 93 | 11.905 | 9.713 | 7.556 | 13.059 |
|  | 11.969 11.776 | . 424 | 7.723 7.517 | 11.785 12.727 | . 90 | . 95 | 97 | 12. | 9.769 | 7.480 | 13.116 |
| $\infty 22$ | 2.075 | 5.279 | 13.789 | 13.7 | 1.0 | . 50 | 50 |  | 9.810 | 7.463 | 13.032 |
| $\infty 11$ | 1.373 | 5.258 | 14.037 | 13.915 | 1.0 | . 00 | . 33 | 1.696 | 11.173 | . 986 |  |
| $\infty 23$ | 2.524 | 11.043 | 9.796 | 12.347 | 1.0 | . 50 | . 67 | 2.793 | 11.285 | 6 | 12.811 |
| $\infty 46$ | 4.097 | 10.848 | 9.402 | 12.092 | 1.0 | 75 | 83 | 4.342 | 11.284 | 9.030 | 12.909 |
| $\infty 69$ | 5.397 | 10.695 | 9.066 | 11.926 | 1.0 | 97 | 98 | 9.565 | 10.450 | 7.845 | 12.737 |
| $1 \times 1$ | 3.547 | 1.530 | 15.798 | 15.233 | . 00 | . | 67 | 15.106 | 2.681 | 11.177 | 14.143 |
| $1 \times 3$ | 15.059 | 2.356 | 11.410 | 13.996 | . 80 | 0 | 93 | 14.749 | 6.100 | 9.017 | 13.875 |
| $2 \infty 6$ | 14.665 | 3.613 | 10.703 | 13.562 | . 90 | 0 | 97 | 13.829 | 7.495 | 8.394 | 13.488 |
| $3 \infty 9$ | 14.301 | 4.662 | 10.123 | 13.220 | . 97 | . 0 | 99 | 12.550 | 8.989 | 7.824 | 12.985 |
| 110 | 7.371 | 20.504 | 1.873 | 17.795 | . 00 | 50 | . 0 | 17.108 | 14.827 | 2.413 | 16.212 |
| 120 | 17.194 | 14.970 | 2.289 | 16.151 | . 80 | 90 | . 0 | 15.525 | 12.011 | 4.817 | 15.117 |
| $24 \infty$ | 16.382 | 14.062 | 3.165 | 15.556 | . 90 | 95 | . 0 | 14.354 | 11.111 | 5.808 | 14.341 |
| $48 \infty$ | 15.069 | 12.717 | 4.498 | 14.616 | . 96 | 98 | 1.0 | 13.084 | 10.405 | 6.702 | 13.520 |
| 1000 | 32.977 | 4.859 | 3.768 | 22.649 | . 50 | 1.0 | 1.0 | 28.947 | 5.770 | 4.505 | 20.897 |
| 2000 | 29.244 | 5.739 | 4.426 | 21.083 | . 75 | 1.0 | 1.0 | 23.877 | 6.929 | 5.421 | 18.503 |
| 40000 | 23.622 | 7.061 | 5.418 | 18.582 | . 90 | 1.0 | 1.0 | 17.715 | 8.370 | 6.515 | 15.362 |
| $8 \infty$ | 17.229 | 8.563 | 6.549 | 15.506 | . 95 | 1.0 | 1.0 | 14.806 | 9.072 | 7.017 | 13.912 |
| $\infty 1 \infty$ | 3.562 | 21.997 | 2.356 | 18.684 | 1.0 | . 50 | 1.0 | 4.647 | 20.467 | 2.806 | 17.875 |
| $\infty 2 \infty$ | 4.322 | 20.861 | 2.652 | 18.098 | 1.0 | . 75 | 1.0 | 6.208 | 18.247 | 3.766 | 16.672 |
| $\infty$ | 5.635 | 18.907 15.990 | 3.517 | 17.072 | 1.0 | . 90 | 1.0 | 8.574 | 14.817 | 5.264 | 14.817 |
| $\infty \times 1$ | 1.952 | 15.990 1.732 | 16.195 | 15.535 15.513 | 1.0 | 1.95 | 1.0 | 9.998 2.789 | 12.687 | 6.209 | 13.754 |
| $\infty \times 2$ | 2.584 | 2.237 | 15.648 | 15.180 | 1.0 | 1.0 | . 75 | 4.124 | 3.787 | 14.101 | 14.255 |
| cos4 | 3.725 | 3.156 | 14.654 | 14.590 | 1.0 | 1.0 | . 90 | 6.563 | 5.981 | 11.825 | 13.118 |
| cose 8 | 5.596 | 4.680 | 13.014 | 13.676 | 1.0 | 1.0 | . | 8.433 | 7.543 | 10.160 | 12.580 |

Table 3b. The model as in table 3a, with the polling order reversed.

| dcp | $\mathrm{E}\left\{\mathrm{W}_{1}\right.$ | $\left\{W_{2}\right\}$ | E \{W | $\sigma\{W$ |  | dcp | $\mathrm{E}\left\{\mathrm{W}_{1}\right.$ | \{ $\left.W_{2}\right\}$ | E\{W | $\sigma\{\mathrm{W}\}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 111 | 2.484 | 5.017 | 13.827 |  | 1.0 | 1.01 .0 | 11.969 | 9.677 | 7.559 | 12.733 |
| 444 | 3.742 | 5.340 | 13.192 | 13.417 | . 75 | . 75.75 | 5.270 | 7.198 | 11.444 | 12.799 |
| 123 | 11.841 | 9.151 | 7.952 | 12.006 | . 00 | . 50.67 | 10.961 | 9.450 | 8.046 | 12.326 |
| 369 | 12.019 | 9.272 | 7.814 | 11.886 | . 90 | . 95.97 | 12.158 | 9.636 | 7.524 | 13.118 |
| $12 \infty$ | 17.228 | 14.957 | 2.286 | 16.154 | . 00 | . 501.0 | 17.084 | 14.827 | 2.421 | 16.211 |
| $1 \times 3$ | 15.068 | 2.358 | 11.405 | 13.995 | . 00 | 1.0 . 67 | 15.204 | 2.662 | 11.158 | 14.150 |
| m23 | 2.526 | 11.018 | 9.812 | 12.346 | 1.0 | . 00.33 | 1.702 | 11.148 | 10.001 | 12.595 |
| 1000 | 33.010 | 4.931 | 3.710 | 22.661 | 1.0 | . 00.00 | 1.408 | 5.223 | 14.048 | 13.919 |
| $\infty 1 \infty$ | 3.367 | 21.989 | 2.218 | 18.674 | . 00 | 1.0 . 00 | 3.671 | 1.481 | 15.789 | 15.231 |
| $\infty \times 1$ | 2.133 | 1.640 | 16.196 | 15.516 | . 00 | . 001.0 | 7.265 | 20.513 | 1.903 | 17.795 |

Table 4a. The waiting time distributions for systems with three queues, equal arrival rates and unequal service rates.

| dcp | $\mathrm{E}\left\{\mathrm{W}_{1}\right\}$ | $\mathrm{E}\left\{\mathrm{W}_{2}\right\}$ |  | E\{ |  | cp |  | , | $\mathrm{E}\left\{\mathrm{W}_{2}\right\}$ | $\mathrm{E}\left\{\mathrm{W}_{3}\right\}$ | E\{W\} |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 111 | 9.09 | 10.147 | 11.203 | 10.1 | . 25 | 25 | . 25 | 9.683 | 16 | 5 |  |
| 222 | 9.316 | 10.225 | 11.078 | 10.206 | . 50 | 50 | . 50 | 10.521 | 10.560 | 10.453 |  |
| 333 | 9.566 9.819 | 10.320 | 10.932 | 10.272 | . 75 | . 75 | . 75 | 11.860 | 10.952 | 9.745 | 10.852 |
| 444 666 | 9.819 10.304 | 10.421 10.616 | 10.780 | 10.340 | . 90 | . 90 | . 90 | 13.152 | 11.343 | 9.054 | 11.183 |
| 888 | 10.304 10.739 | 10.616 10.803 | 10.488 10.208 |  | . 95 | 95 | 95 | 13.666 | 11.530 | 8.758 | 11.318 |
| 123 | 32.432 | 8.79 | 4.326 |  | . 99 | 99 | 99 | 13.87 | 11.728 | 8.555 |  |
| 246 | 31.837 | 8.825 | 4.465 |  | . 00 | . 50 | . 67 | 31.144 | 8.673 | 4.837 | 14.884 |
| 369 | 30.905 | 9.044 | 4.671 | 14.859 14.869 |  |  |  | 28.719 | 9.022 | 5.412 | 14.384 |
| 4812 | 29.961 | 9.159 | 4.907 | 14.676 | . 90 | 95 | . 97 |  | 10.550 | 082 |  |
| -0000 | 13.755 | 11.786 | 8.558 | 11.366 | . 96 |  | . 99 | 17.130 | 11.212 | 7.815 | 12.052 |
| 211 | 2.525 | 11.406 | 12.554 | 8.828 | . 50 | 00 | . 0 | 3.060 | 11.304 | 12.444 | 8.936 |
| 811 | 1.664 | 11.574 | 12.729 | 8.656 | . 9 | 00 | . 00 | 1.843 | 11.539 | 12.693 | 8.692 |
| $\infty 11$ | 1.660 | 11.575 | 12.730 | 8.655 | 1.0 | 50 | . 50 | 2.835 | 12.143 | 11.959 | 8.979 |
| $\infty 22$ | 2.581 | 11.527 | 12.455 | 8.854 | 1.0 | . 90 | . 90 | 7.403 | 12.805 | 9.996 | 68 |
| $\infty 44$ | 4.208 | 11.522 | 11.916 | 9.215 | 1.0 | . 95 | . 95 | 9.526 | 12.639 | 9.398 | 1 |
| 121 | 12.262 | 3.084 | 14.857 | 10.067 | . 00 | . 50 | . 00 | 12.075 | 3.532 | 14.620 | 6 |
| 1 | 12.862 | 1.800 | 15.513 | 10.058 | . 00 | . 90 | . 00 | 12.772 | . 00 | 15.408 | 1 |
| $1 \infty 1$ | 12.877 | 1.768 | 15.529 | 10.058 | . 50 | . 0 | . 50 | 14. | 2.797 | 14.321 |  |
| $2 \infty 2$ | 12.824 | 2.519 | 15.046 | 10.130 | . 90 | . 0 | . 90 | 16.4 | 6.809 | 10.970 |  |
| 404 | 12.856 | 3.864 | 14.139 | 10.286 | . 95 | . 0 | . 95 | 15.959 | 8.588 | 9.955 |  |
| 112 | 16.044 | 17.466 | 4.008 | 12.506 | . 00 | . 00 | . 50 | 15.744 | 17.132 | 4.331 | 12.402 |
| 118 | 17.786 | 19.254 | 2.235 | 13.092 | . 00 | 00 | 90 | 17.609 | 19.067 | 2.419 | 13.032 |
| 11ヵ | 17.900 | 19.369 | 2.121 | 13.130 | . 50 | 50 | . 0 | 18.251 | 18.208 | 2.778 | 13.079 |
| $22 \infty$ | 17.516 | 18.809 | 2.622 | 12.982 | . 90 | 90 | 1.0 | 17.770 | 14.665 | 5.300 | 12.578 |
| $12 \infty$ | 16.864 | 17.818 | 3.500 | 12.727 | . 95 | . 95 | 1.0 | 16.741 | 13.480 | 6.433 | 18 |
| $12 \infty$ | 33.779 | 11.002 | 2.406 | 15.729 | . 00 | . 50 | 1.0 | 33.300 | 11.068 | 2.522 | 30 |
| 2000 | 40.320 | 5.445 | 3.930 | 16.565 | . 50 | 1.0 | 1.0 | 37.775 | 6.025 | 4.392 | 64 |
| 20 | 38.112 34.240 | 5.971 6.896 | 4.315 | 16.133 | . 75 | 1.0 | 1.0 | 33.782 | 6.943 | 5.111 | 15.278 |
| 80000 | 34.240 28.176 | . 896 | 4.990 | 15.375 | . 90 | 1.0 | 1.0 | 26.636 | 8.611 | 6.381 | 13.876 |
| $\infty 1 \infty$ | 4.277 | 25.762 | 2.400 | 10.81 | 1.95 | 1.0 | 1.0 | 21.403 | 9.860 | 7.293 |  |
| $\infty$ 2m | 5.032 | 24.639 | 2.897 | 10.856 | 1.0 | . 75 | 1.0 | 5.374 7.006 | 24.221 21.911 | 3.061 4.057 | 10.991 |
| $\infty 40$ | 6.352 | 22.680 | 3.763 | 10.932 | 1.0 | . 90 | 1.0 | 9.658 | 18.095 | 5.718 | 11.157 |
| $\infty 8 \infty$ | 8.403 | 19.639 | 5.103 | 11.050 | 1.0 | . 95 | 1.0 | 11.397 | 15.520 | 6.854 | 11.257 |
| $\infty \times 1$ | 2.327 | 2.068 | 18.845 | 7.747 | 1.0 | 1.0 | . 50 | 3.524 | 3.227 | 17.674 | 8.142 |
| $\infty \times 2$ | 3.243 | 2.814 | 18.043 | 8.033 | 1.0 | 1.0 | . 75 | 5.340 | 4.955 | 15.917 | 8.737 |
| $\infty \times 04$ | 4.853 | 4.145 | 16.619 | 8.539 | 1.0 | 1.0 | . 90 | 8.370 | 7.712 | 13.069 | 9.717 |
| $\infty \times 08$ | 7.357 | 6.252 | 14.380 | 9.330 | 1.0 | 1.0 |  | 10.457 | 9.467 | 11.203 | 10.376 |

Table 4b. The model as in table 4a, with the polling order reversed.

| dcp | $\mathrm{E}\left\{\mathrm{W}_{1}\right\}$ | $\mathrm{E}\left\{\mathrm{W}_{2}\right\}$ | $\mathrm{E}\left\{\mathrm{W}_{3}\right\}$ | $\mathrm{E}\{\mathrm{W}\}$ | dcp | $\mathrm{E}\left\{\mathrm{W}_{1}\right\}$ | $\mathrm{E}\left\{\mathrm{W}_{2}\right\}$ | $\mathrm{E}\left\{\mathrm{W}_{3}\right\}$ | $\mathrm{E}\{\mathrm{W}\}$ |  |
| :--- | ---: | ---: | ---: | ---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 111 | 9.130 | 10.108 | 11.218 | 10.152 | 1.0 | 1.0 | 1.0 | 14.397 | 11.257 | 8.696 |
| 11.450 |  |  |  |  |  |  |  |  |  |  |
| 444 | 10.006 | 10.218 | 10.852 | 10.359 | .75 | .75 | .75 | 11.997 | 10.817 | 9.790 |
| 10.868 |  |  |  |  |  |  |  |  |  |  |
| 123 | 32.491 | 8.792 | 4.312 | 15.196 | .00 | .50 | .67 | 31.196 | 8.656 | 4.831 |
| 369 | 31.172 | 8.891 | 4.685 | 14.911 | .90 | .95 | .97 | 20.984 | 10.355 | 7.102 |
| 12.894 |  |  |  |  |  |  |  |  |  |  |
| $10 \infty$ | 40.365 | 5.465 | 3.902 | 16.577 | 1.0 | .00 | .00 | 1.674 | 11.515 | 12.765 |
| $\infty 1 \infty$ | 4.102 | 25.727 | 2.481 | 10.770 | .00 | 1.0 | .00 | 12.943 | 1.748 | 15.520 |
| $\infty 00$ | 10.070 |  |  |  |  |  |  |  |  |  |
|  | 2.606 | 1.913 | 18.856 | 7.792 | .00 | .00 | 1.0 | 17.951 | 19.329 | 2.130 |

Table 5a. The influence of one relatively heavily loaded queue, $\rho=0.75$.


Table 5b. The influence of one relatively heavily loaded queue, $\rho=0.75$.

|  | s $=4, \quad$ limited |  |  | dcp | ( $\mathrm{s}=4, \quad$ Bernoulli |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| dcp | $\mathrm{E}\left\{\mathrm{W}_{1}\right\} \mathrm{E}\left\{\mathrm{W}_{2}\right\} \mathrm{E}\left\{\mathrm{W}_{4}\right.$ | $E\{W\}$ | $\sigma\{\mathrm{W}\}$ |  | $\mathrm{E}\left\{\mathrm{W}_{1}\right\} \mathrm{E}\left\{\mathrm{W}_{2}\right\}$ | E $\left\{\mathrm{W}_{4}\right\}$ | $E\{W\}$ | $\sigma\{W\}$ |
| 11 | 3.0221 .4471 .495 | 2.091 | 3.273 | . 25.00 | 2.7341 .828 | 1.882 | 2.206 |  |
| 21 | $2.2402 .485 \quad 2.542$ | 2.404 | 3.627 | . 50.00 | 2.266 2.447 | 2.510 | 2.206 2.394 | 3.319 3.677 |
| $\begin{array}{ll}3 & 1 \\ 4 & 1\end{array}$ | $\begin{array}{llll}1.754 & 3.123 & 3.199 \\ 1.521 & 3.429 & 3.515\end{array}$ | 2.598 | 4.286 | . 67.00 | 1.8682 .973 | 3.046 | 2.553 | 4.222 |
| 4 6 6 | $\begin{array}{llll}1.521 & 3.429 & 3.515 \\ 1.325 & 3.687 & 3.782\end{array}$ | 2.691 | 4.668 | . 75.00 | 1.6723 .232 | 3.311 | 2.631 | 4.536 |
|  | $\begin{array}{llll}1.325 & 3.687 & 3.782 \\ 1.170 & 3.891 & 3.990\end{array}$ | 2.770 2.832 | 5.011 | . 90.00 | 1.3533 .651 | 3.742 | 2.759 | 5.056 |
| $\infty 2$ | $1.3553 .591 \quad 3.799$ | 2.758 | 5.112 | [99 1.00 | 1.187 <br> 1.328 <br> 1.868 <br> 1.650 | 3.967 3.810 | 2.825 2.769 | 5.321 5.212 |
| $\infty 4$ | $\begin{array}{lllll}1.573 & 3.235 & 3.580\end{array}$ | 2.671 | 4.795 | 1.0 | 1.4793 .418 | 3.642 | 2.708 |  |
| 4 6 6 | $\begin{array}{llll}2.138 & 2.582 & 2.721 \\ 2.048 & 2.664 & 2.877\end{array}$ | 2.445 | 3.676 | . 75.50 | 2.0972 .648 | 2.761 | 2.461 |  |
| 22 | $\begin{array}{llll}2.048 & 2.664 & 2.877 \\ 2.894 & 1.598 & 1.687\end{array}$ | 2.481 | 3.740 | . 95.90 | 1.8622 .888 | 3.153 | 2.555 |  |
|  | 2.7541 .7671 .894 | 2.198 | 3.265 | . 50.50 | $\begin{array}{lll}2.703 & 1.852 \\ 2.388 & 2.243\end{array}$ | 1.941 | 2.219 | 3. |
|  | 2.6211 .9302 .090 | 2.252 | 3.299 | . 90.90 | 2.079 2.612 |  | 2.469 |  |
|  | 1.7932 .9263 .312 | 2.583 | 4.412 | . 95.95 | 1.9442 .765 | 3.063 | 2.523 | 4. |
| $\begin{array}{ll}1 \\ 1 & 2 \\ 1\end{array}$ | 3.3081 .0491 .133 | 1.977 | 3.344 | . 00.50 | 3.2291 .163 | 1.229 | 2.009 | 3. |
| $\begin{aligned} & 1 \infty \\ & 2 \infty \end{aligned}$ | 3.3890 .9331 .036 | 1.945 | 3.389 | . 00.75 | 3.3141 .043 | 1.124 | 1.975 | 3. |
|  | 3.1191 .2891 .401 | 2.052 | 3.306 | . 501.0 | 3.0561 .347 | 1.511 | 2.0 |  |
| 6 | 2.7131 .8101 .964 | 2.215 | 3.300 | . 751.0 | 2.6741 .824 | 2.059 | 2.2 | 3.41 |
|  | $2.434 \quad 2.158 \quad 2.361$ | 2.326 | 3.411 | .901 .0 | 2.2412 .365 | 2.676 | 2.404 | 3. |
|  | s $=5$, limited |  |  | dcp | s $=5$, Bernoull |  |  |  |
| dcp | $\mathrm{E}\left\{\mathrm{W}_{1}\right\} \mathrm{E}\left\{\mathrm{W}_{2}\right\} \quad \mathrm{E}\left\{\mathrm{W}_{5}\right\}$ | $\mathrm{E}\{\mathrm{W}$ | $\sigma\{\mathrm{W}\}$ |  | $\mathrm{E}\left\{\mathrm{W}_{1}\right\} \quad \mathrm{E}\left\{\mathrm{W}_{2}\right\}$ | $E\left\{W_{5}\right\}$ | E \{ W$\}$ | $\sigma\{\mathrm{W}\}$ |
| 11 | 2.9911 .4811 .540 | 2.00 | 3.18 | 25.00 | 2.6571 .811 | 1.877 | 2.114 |  |
| 21 | 2.1032 .3612 .433 | 2.299 | 3.55 | . 50.00 | 2.1592 .302 | 2.380 | 2.280 | 3.57 |
| 31 | 1.6222 .8332 .935 | 2.461 | 4.09 | . 67.00 | 1.7682 .688 | 2.778 | 11 | 4.00 |
| 41 | 1.4003 .0573 .160 | 2.532 | 4.37 | . 75.00 | 1.5792 .873 | 2.971 | 2.474 | 4. |
| 51 | 1.2883 .1603 .250 | 2.574 | 4.50 | . 90.00 | 1.2773 .168 | 3.280 | 2.574 | 4.63 |
|  | 1.1043 .3363 .457 | 2.632 | 4.85 | . 95.00 | 1.1883 .255 | 3.371 | 2.604 | 4.74 |
|  | $\begin{array}{llll}1.325 & 3.065 & 3.296 \\ 1.994 & 2.423 & 2.599\end{array}$ | 2.559 | 4.60 | 1.0 . 50 | 1.2773 .133 | 3.317 | 2.574 | 4. |
|  | $\begin{array}{llll}1.994 & 2.423 & 2.599 \\ 2.814 & 1.629 & 1.736\end{array}$ | 2.335 | 3.56 | . 75.50 | 1.9932 .442 | 2.576 | 2.336 | 3.79 |
|  | 1.7372 .5942 .962 | 2.421 | 4.17 | . 50.50 | 2.6211 .828 | 1.934 | 2.126 | 3. |
| 12 | 3.3471 .1061 .206 | 1.884 | 3.25 | . 75.75 | 2.2852 .134 | 2.305 | 2.238 | 3.5 |
|  | 3.4451 .0021 .119 | 1.852 | 3.29 | . 00.75 | 3.2431 .220 <br> 3.349 | 1.299 1.201 | 1.919 1.884 | 3.23 |
|  | 3.0851 .3571 .486 | 1.972 | 3.20 | . 501.0 | 3.0311 .386 | 1.568 | 1.889 | 3.26 |
|  | 2.5821 .8382 .018 | 2.139 | 3.22 | .751 .0 | 2.5911 .796 |  | 2.136 | 31 |
|  | s $s=6$, limited |  |  | dcp | s = 6, Bernoulli |  |  |  |
| dcp | $E\left\{W_{1}\right\} \quad E\left\{W_{2}\right\} \quad \mathrm{E}\left\{\mathrm{W}_{6}\right.$ | E\{W\} | $\sigma\{\mathrm{W}\}$ |  | $E\left\{W_{1}\right\} E\left\{W_{2}\right\}$ | $\mathrm{E}\left\{\mathrm{~W}_{6}\right\}$ | $E\{W\}$ | \{ |
| 11 | 2.9551 .5071 .573 | 1.942 | 3.10 | 25.00 | 2.5891 .796 | 1.871 |  |  |
| 21 | $1.9962 .263 \quad 2.345$ | 2.215 | 3.47 | . 50.00 | 2.0802 .194 | 2.281 | 2.191 | 3.47 |
| 41 | $\begin{array}{lllll}1.327 & 2.751 & 2.919\end{array}$ | 2.398 | 4.11 | . 75.00 | 1.5252 .628 | 2.735 | 2.350 | 4.01 |
|  | $\begin{array}{llll}1.074 & 2.975 & 3.108 \\ 1.323 & 2.732 & 2.971\end{array}$ | 2.478 | 4.49 | . 95.00 | 1.1532 .915 | 3.043 | 2.456 | 4.41 |
|  | $\begin{array}{llll}1.323 & 2.732 & 2.971 \\ 2.746 & 1.651 & 1.764\end{array}$ | 2.410 2.001 | 4.25 | $\begin{array}{rrr}1.0 & 50\end{array}$ | 1.2562 .802 | 2.996 | 2.427 | 4.38 |
|  | $\begin{array}{llll}1.746 & 1.651 & 1.764 \\ 1.700 & 2.387 & 2.729\end{array}$ | 2.001 | 3.07 | . 50.50 | 2.5491 .808 | 1.924 | 2.058 | 3.22 |
| 42 | 1.8882 .3012 .489 | 2.246 | 3.80 | . 75.75 | 2.2052 .054 | 2.232 | 2.156 | 3.42 |
| 12 | 3.3651 .1581 .265 | 1.824 |  | .75 .00 .00 .50 | $\begin{array}{ll}1.921 & 2.294 \\ 3.240 & 1.268\end{array}$ | 2.438 | 2.237 | 3.65 |
|  | 3.4791 .0631 .187 | 1.792 | 3.21 | . 00 . 75 | 3.2401 .268 | 1.354 | 1.860 | 3.15 |
|  | 3.0331 .4131 .550 | 1.919 | 3.12 | . 501.0 | $\begin{array}{ll}3.365 & 1.162 \\ 2.996 & 1.422\end{array}$ | 1.264 | 1.824 | 3.18 3.15 |
| $4 \infty$ | 2.4522 .5152 .844 | 2.085 | 3.15 | .751 .0 | 2.5181 .777 | 2.02 | 2.066 | 3.15 3.23 |

Table 6. The influence of one relatively heavily loaded queue, $p=0.95$.


## Appendix

The examples described in section 5 will be further elaborated in this appendix. The heavy traffic behaviour of the moments of the waiting time distributions for the model of table 6 is shown in table A. 1 as function of $q_{1}$ for $q_{2}=0.00$. Note the strong sensitivity, especially of $E\left\{W_{2}\right\}$, with respect to the parameter $q_{1}$ in the neighborhood of $q_{1}=0.50$. For $q_{1}=0.50$ the queues are balanced in the sense that the length of both queues tends to infinity in heavy traffic. Table A. 2 contains estimations of heavy traffic limits for the same model, but for Bernoulli schedules with $a_{1}\left(1-q_{1}\right)=a_{2}\left(1-q_{2}\right)$, the case which has not been covered by the conjecture of property 2. Although limited service and Bernoulli schedules agree in the sense of property 1 , the actual values of the limits $\omega_{j}$, $j=0,1, \ldots, s$, behave quite differently for the two service disciplines. For limited service with $K_{1}=2 \mathrm{~K}_{2}$ we find the following estimates (these limits seem to vary hardly with $K_{2}$, as far as they can be determined with sufficient accuracy) :

Table A.1. Model of tables 5,6, Bernoulli schedules.

| dcp | $s=2, \quad p=0.98$ |  |  | $\sigma\{W\}$ | $s=2, \quad \rho=0.99$ |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| . 00.00 | 54.48 | 2.59 | 37.18 | 52.0 | 110.71 | 2.67 | 74.69 | 105.0 |
| . 25.00 | 53.91 | 4.87 | 37.56 | 51.3 | 110.10 | 5.11 | 75.10 | 104.3 |
| . 45.00 | 50.61 | 18.05 | 39.76 | 48.5 | 105.98 | 21.58 | 77.83 | 100.1 |
| . 49.00 | 45.59 | 38.14 | 43.11 | 48.8 | 96.11 | 60.85 | 84.43 | 95.7 |
| . 50.00 | 42.65 | 49.91 | 45.07 | 52.5 | 86.11 | 101.06 | 91.09 | 105.3 |
| . 51.00 | 38.82 | 65.21 | 47.62 | 61.3 | 70.35 | 164.14 | 101.55 | 147.1 |
| . 55.00 | 25.51 | 118.43 | 56.49 | 104.0 | 34.99 | 305.20 | 125.16 | 259.2 |
| . 67.00 | 11.73 | 173.59 | 65.68 | 143.6 | 13.38 | 391.98 | 139.58 | 311.0 |
| . 75.00 | 8.25 | 187.49 | 68.00 | 151.7 | 9.05 | 409.28 | 142.46 | 319.8 |
| . 90.00 | 5.20 | 199.69 | 70.03 | 158.1 | 5.52 | 423.41 | 144.82 | 326.0 |
| 1.0 .00 | 4.08 | 204.17 | 70.78 | 160.3 | 4.28 | 428.37 | 145.64 | 328.1 |

Table A.2. Model of tables 5,6: estimations of heavy traffic limits.

| dcp | s $=2$, Bernoulli |  |  |  | $\mathbf{s}=3$, Bernoulli |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $\omega_{1}$ | $\omega_{2}$ | $\omega_{0}$ | ${ }^{0} 0$ | $\omega_{1}$ | $\omega_{2}$ | $\omega_{3}$ | $\omega_{0}$ | $\zeta_{0}$ |
| .50 .00 | 0.869 | 1.023 | 0.921 | 1.06 | 0.788 | 0.925 | 0.924 | 0.856 | 1.05 |
| . 75.50 | 0.815 | 1.239 | 0.957 | 1.18 | 0.715 | 1.070 | 1.070 | 0.893 | 1.18 |
| .90 .80 | 0.736 | 1.558 | 1.010 | 1.40 | 0.624 | 1.253 | 1.252 | 0.938 | 1.36 |
| . 95.90 | 0.679 | 1.783 | 1.047 | 1.57 | 0.571 | 1.359 | 1.357 | 0.965 | 1.48 |
| .98 .96 | 0.623 | 2.008 | 1.085 | 1.75 | 0.528 | 1.444 | 1.442 | 0.986 | 1.58 |
| .99 .98 | 0.596 | 2.115 | 1.102 | 1.83 | 0.518 | 1.463 | 1.465 | 0.991 | 1.61 |
| 1.01 .0 | 0.563 | 2.250 | 1.125 | 1.95 | 0.556 | 1.389 | 1.389 | 0.972 | 1.43 |

$$
\begin{array}{lll}
s=2: & \omega_{1}=0.889, & \omega_{2}=0.946,
\end{array} \quad \omega_{0}=0.908, \quad \zeta_{0}=1.02 ;
$$

The limits in table A. 2 have been determined by means of extrapolation from values of the performance measures for values of $\rho$ close to 1 . It is not clear whether differences between $\omega_{2}$ and $\omega_{3}$ (case $s=3$ ) exist or stem from inaccuracies of the computations. In the case of exhaustive service these limits are equal by theorem 2. Above we have used, in addition to (4.5), the following notation:

$$
\begin{equation*}
\zeta_{0}:=\lim _{\rho \uparrow 1}(1-\rho) \sigma\{W\} . \tag{A.1}
\end{equation*}
$$

Next, we will consider the result of theorem 1a in more detail. Firstly, it is readily verified with the aid of (2.4) that, for all vectors $\bar{q}$,

$$
\begin{equation*}
\sum_{j=1}^{s} \eta_{j} \Psi(j, \bar{q})=0 ; \tag{A.2}
\end{equation*}
$$

hence, by the conservation law, cf. (4.3), it must hold that

$$
\begin{equation*}
\frac{\beta_{2}}{2 \beta_{1}} \sum_{j=1}^{s} \eta_{j}^{2}+\sum_{j=1}^{s} \eta_{j} \Xi(j)=\frac{\beta_{2}}{2 \beta_{1}} \tag{A.3}
\end{equation*}
$$

For the job weighted mean waiting time, cf. (4.2), it is found that

$$
\begin{align*}
E\{W\}= & \rho \frac{\beta_{2}}{2 \beta_{1}}+\rho^{2}\left[\frac{1}{2} \beta_{2} \sum_{j=1}^{S} a_{j} \eta_{j}+\beta_{1} \sum_{j=1}^{s} a_{j} \equiv(j)+2 \sum_{j=1}^{S} q_{j} n_{j}^{2}\left(\frac{1}{\mu_{j}}-\beta_{1}\right)\right] \\
& +O\left(\rho^{3}\right), \quad \text { as } \rho \downarrow 0 . \tag{A.4}
\end{align*}
$$

This expansion indicates that $E\{W\}$ is minimal in light traffic when $q_{j}=1$ if $1 / \mu_{j}<\beta_{1}$, and $q_{j}=0$ if $1 / \mu_{j}>\beta_{1}$, for $j=1, \ldots$ s (cf. the end of section 6). If all service rates are equal (i.e. $\mu_{j}=\mu, j=1, \ldots, s$ ), then the quantities appearing in the power-series expansion (4.4) become, for $j=1, \ldots, s$,

$$
\begin{align*}
& \equiv(j)=\frac{2}{\mu} \sum_{i=1}^{s-1} \eta_{j+i} \sum_{\nu=0}^{i-1} \eta_{j+\nu}=\frac{2}{\mu} \sum_{i=2}^{s} \eta_{i} \sum_{\nu=1}^{i-1} \eta_{\nu}, \quad \text { (independent of } j \text { ), } \\
& \Psi(j, \bar{q})=\frac{2}{\mu}\left[\sum_{i=1}^{s} q_{i} \eta_{i}^{2}-q_{j} \eta_{j}\right], \tag{A.5}
\end{align*} \quad \sum_{j=1}^{s} a_{j} \Psi(j, \bar{q})=0 . \quad \text { (A.5 }, ~ l
$$

If, moreover, the arrival rates are equal (i.e. $\eta_{j}=1 / s, j=1, \ldots, s$ ), then

$$
\equiv(j)=\frac{s-1}{s} \frac{1}{\mu}, \quad \Psi(j, \bar{q})=\frac{2}{s \mu}\left[\frac{1}{s} \sum_{i=1}^{s} q_{i}-q_{j}\right], \quad j=1, \ldots, s ;
$$

note that $\Psi(j, \bar{q})$ vanishes when also $q_{j}=q$ for all $j, j=1, \ldots, s$. For models with $a_{j}=a_{2}, \mu_{j}=\mu_{2}, j=2, \ldots, s$, (cf. tables 5 and 6) we have

$$
\begin{aligned}
& \frac{\beta_{2}}{2 \beta_{1}}=\frac{\eta_{1}}{\mu_{1}}+(s-1) \frac{\eta_{2}}{\mu_{2}} ; \quad \equiv(1)=2(s-1) n_{2}\left[\frac{n_{1}}{\mu_{1}}+\frac{s-2}{2} \frac{\eta_{2}}{\mu_{2}}\right] ; \\
& \equiv(j)=2 \eta_{1}(s+1-j) \frac{\eta_{2}}{\mu_{2}}+2 \eta_{2}(j-2) \frac{\eta_{1}}{\mu_{1}}+\eta_{2}(s-1)(s-2) \frac{\eta_{2}}{\mu_{2}}, \quad j=2, \ldots, s ; \\
& \Psi(j, \bar{q})=2\left[q_{1} \frac{n_{1}^{2}}{\mu_{1}}+\frac{n_{2}^{2}}{\mu_{2}} \sum_{i=2}^{s} q_{i}-q_{j} \frac{\eta_{j}}{\mu_{j}}\right], \quad j=1, \ldots, s .
\end{aligned}
$$

Note that $\equiv(j)$ is increasing with $j, j=2, \ldots, s$, if and only if $1 / \mu_{1}>1 / \mu_{2}$. For the specific models considered in the tables 5 and 6 these quantities become:

$$
\begin{aligned}
& \beta_{1}=\frac{s+3}{2(s+1)}, \quad \frac{\beta_{2}}{2 \beta_{1}}=\frac{s+7}{2(s+3)}, \quad \eta_{1}=\frac{4}{s+3}, \quad \eta_{j}=\frac{1}{s+3}, \quad j=2, \ldots, s ; \\
& \equiv(1)=\frac{(s-1)(s+14)}{2(s+3)^{2}}, \quad \equiv(j)=\frac{(s-1)(s-2)}{2(s+3)^{2}}+\frac{4(s+j-3)}{(s+3)^{2}}, \quad j=2, \ldots, s ; \\
& \Psi(1, \bar{q})=\frac{1}{(s+3)^{2}}\left[\sum_{i=2}^{s} q_{i}-8(s-1) q_{1}\right]=\frac{s-1}{(s+3)^{2}}\left(q_{2}-8 q_{1}\right) ; \\
& \Psi(j, \bar{q})=\frac{1}{(s+3)^{2}}\left[32 q_{1}+\sum_{i=2}^{s} q_{i}-(s+3) q_{j}\right]=\frac{4}{(s+3)^{2}}\left(8 q_{1}-q_{2}\right), \quad j=2, \ldots, s .
\end{aligned}
$$

For general models with two queues these quantities become:

$$
\begin{aligned}
& \Xi(j)=2 \eta_{1} \eta_{2} / \mu_{j}, \quad j=1,2 ; \\
& \Psi(1, \bar{q})=2 \eta_{2}\left[q_{2} \frac{\eta_{2}}{\mu_{2}}-q_{1} \frac{\eta_{1}}{\mu_{1}}\right], \quad \Psi(2, \bar{q})=2 \eta_{1}\left[q_{1} \frac{\eta_{1}}{\mu_{1}}-q_{2} \frac{\eta_{2}}{\mu_{2}}\right] .
\end{aligned}
$$

From these relations it readily follows that $E\left\{W_{1}\right\}>E\left\{W_{2}\right\}$ for small values of $\rho$ if $\left(1-2 q_{1}\right) \eta_{1} / \mu_{1}>\left(1-2 q_{2}\right) \eta_{2} / \mu_{2}$. On the other hand, property 1 implies that $E\left\{W_{1}\right\}>E\left\{W_{2}\right\}$ for values of $\rho$ close to 1 if $a_{1}\left(1-q_{1}\right)>$ $a_{2}\left(1-q_{2}\right)$. For the model of tables 5 and 6 , with 2 queues, the above implies that there is a different ordering of $E\left\{W_{1}\right\}$ and $E\left\{W_{2}\right\}$ in light and heavy traffic if $7+2 q_{2}<16 q_{1}<8+8 q_{2}$. This feature is illustrated in table A. 3 where the mean waiting times are shown for several values of the load $\rho$ for service disciplines with $q_{1}=q_{2}=q$ (respectively $K_{1}=K_{2}=K$ ). Note in this table also the close resemblence of the mean waiting times in the cases $K=16$ and $q=1.0(K=\infty)$ for $p \leq 0.5$ ( $c f$. the observation below theorem $1)$, and the important differences for larger values of $\rho$. This rather abrupt change in the behaviour of the performance measures causes slow convergence of the power-series for larger values of $\rho$ for such models. At this point it is also important to note the discontinuity in the heavy traffic limits of the mean waiting times, for instance when $q_{1}=q_{2}=q$ and $q$ approaches 1 : for $q<1$ these limits are determined by corollary 1 , while for $q=1$ they are given by theorem 2 (see also table A.3). The quantities appearing in theorem 1 become for the model of table 3:

$$
\begin{array}{ll}
\beta_{1}=1, \quad \frac{\beta_{2}}{2 \beta_{1}}=1, \quad \eta_{1}=\frac{1}{6}, & \eta_{2}=\frac{1}{3}, \quad \eta_{3}=\frac{1}{2} ; \\
\Psi(1, \bar{q})=\frac{2}{9} q_{2}+\frac{1}{2} q_{3}-\frac{5}{18} q_{1}, & \Psi(2, \bar{q})=\frac{1}{18} q_{1}+\frac{1}{2} q_{3}-\frac{4}{9} q_{2}, \\
\Psi(3, \bar{q})=\frac{1}{18} q_{1}+\frac{2}{9} q_{2}-\frac{1}{2} q_{3} ; & \equiv(j)=\frac{11}{18}, \quad j=1,2,3 .
\end{array}
$$

The values of $\Xi(j)$ are independent of $j$ and of the polling order for models with equal service rates, cf. (A.5). However, higher order terms of the power-series expansions of the mean waiting times do depend on the polling order; see the differences in the tables $3 a$ and $3 b$.
The quantities appearing in theorem 1 become for the model of table 4:

$$
\begin{aligned}
& \beta_{1}=1, \quad \frac{\beta_{2}}{2 \beta_{1}}=\frac{7}{6}, \quad \eta_{1}=\frac{1}{6}, \quad \eta_{2}=\frac{1}{3}, \quad \eta_{3}=\frac{1}{2} ; \\
& \Xi(1)=\frac{17}{36}, \quad \Xi(2)=\frac{25}{36}, \quad \Xi(3)=\frac{29}{36}, \quad \sum_{j=1}^{3} a_{j} \Xi(j)=\frac{71}{108} ; \\
& \Psi(1, \bar{q})=\frac{2}{9} q_{2}+\frac{3}{4} q_{3}-\frac{5}{36} q_{1}, \quad \Psi(2, \bar{q})=\frac{1}{36} q_{1}+\frac{3}{4} q_{3}-\frac{4}{9} q_{2}, \\
& \Psi(3, \bar{q})=\frac{1}{36} q_{1}+\frac{2}{9} q_{2}-\frac{3}{4} q_{3}, \quad \sum_{j=1}^{3} a_{j} \Psi(j, \bar{q})=\frac{1}{4} q_{3}-\frac{1}{36} q_{1} .
\end{aligned}
$$

When the polling order is reversed in this model then we have:

$$
\Xi(1)=\frac{23}{36}, \quad \Xi(2)=\frac{19}{36}, \quad \Xi(3)=\frac{31}{36}, \quad \sum_{j=1}^{3} a_{j} \Xi(j)=\frac{73}{108} .
$$

Tables A. 4 and A. 5 contain estimations for heavy traffic limits of waiting time characteristics for the models of table 3 and table 4 respectively. These tables should be read as follows. An entry without brackets stands for an estimation for $\omega_{j}, j=1,2,3$; an entry between brackets indicates that $\omega_{j}=0$, cf. property 1 , and stands for an estimation for $x_{j}, j=1,2,3$, cf. (4.6). For the model of table 3 we have $\omega_{0}=1.00$, independent of the service discipline.

Table A.3. Model of tables 5,$6 ; 2$ queues, $q_{1}=q_{2}=q$ respectively $K_{1}=K_{2}=K$.

| $q=.00$ | $p=0.1$ | $\rho=0.5$ | $\rho=0.7$ | $\rho=0.8$ | $p=0.9$ | $p=0.92$ | $p=0.94$ | $p=0.96$ | $\rho=0.98$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $E\left\{W_{1}\right\}$ | . 1014 | 0.9606 | 2.3313 | 4.1119 | 9.6107 | 12.39 | 17.05 | 27.95 | 8 |
| $E\left\{W_{2}\right\}$ | . 0944 | 0.6576 | 1.1747 | 1.5525 | 2.0574 | 2.18 | 2.31 | 2.45 | 2.59 |
| $q=.50$ | $p=0.1$ | $\rho=0.5$ | $p=0.7$ | $p=0.8$ | $p=0.9$ | $p=0.92$ | $p=0.94$ | $p=0.96$ | $p=0.98$ |
| $\mathrm{E}\left\{\mathrm{W}_{1}\right\}$ | . 1001 | 0.9192 | 2.2201 | 3.9309 | 9.3125 | 12.06 | 16.68 | 25.98 | 3 |
| $\mathrm{E}\left\{\mathrm{W}_{2}\right\}$ | . 0998 | 0.8232 | 1.6 | 2.276 | 3.249 | 3.50 | 3.77 | 4.07 | 4.39 |
| $q=.75$ | $p=0.1$ | $p=0.5$ | $p=0.7$ | $p=0.8$ | $\rho=0.9$ | $p=0.92$ | $p=0.94$ | $p=0.96$ | $\rho=0.98$ |
|  | . 0993 | 0.8 | 2.0902 | 3.6 | 8.8413 | 11.52 | 16.05 | 25.25 | 53.17 |
| $\mathrm{E}\left\{\mathrm{W}_{2}\right\}$ | . 1030 | 0.9818 | 2.139 | 3.237 | 5.1348 | 5.68 | 6.29 | 7.00 | 7.81 |
| $q=.90$ | $p=0.1$ | $p=0$ | $\rho=0.7$ | $\rho=0.8$ | $\rho=0.9$ | $p=0.92$ | $p=0.94$ | $p=0.96$ | $p=0.98$ |
| $E\left\{W_{1}\right\}$ | . 098 | 0.8 | 1.9256 | 3.337 | 7.9 | 10.42 | 14.68 | 23.50 | 50.90 |
| $\mathrm{E}\left\{\mathrm{W}_{2}\right\}$ | . 105 | 1.1 | 2.7 | 4. | 8.6 | 10.05 | 11.79 | 14.01 | 16.91 |
| $\mathrm{q}=.95$ | $\rho=0.1$ | $p=0$ | $\rho=0.7$ | $\rho=0.8$ | $\rho=0.9$ | $p=0.92$ | $p=0.94$ | $p=0.96$ | $p=0.98$ |
| $\mathrm{E}\left\{\mathrm{W}_{1}\right\}$ | . 0 | 0.8 | 1 | 3.105 | 7.2217 | 9.43 | 13.30 | 21.50 | 8 |
| $E\left\{W_{2}\right\}$ | . 1059 | 1.2 | 3.181 | 5. | 11. | 14.01 | 17.29 | 22.00 | 29.14 |
| $\mathrm{q}=.99$ | $p=0.1$ | $p=0.5$ | $p=0.7$ | $\rho=0.8$ | $p=0.9$ | $p=0.92$ | $p=0.94$ | $p=0.96$ | $p=0.98$ |
| $E\left\{W_{1}\right\}$ | . 0 | 0 | 1.7168 | 2.8031 | 6.0571 | 7.73 | 10.61 | 16 | 37.15 |
| $\mathrm{E}\left\{\mathrm{W}_{2}\right\}$ | . 1066 | 1.30 | 3 | 6. | 16.2 | 20.83 | 28.07 | 41.25 | 71.89 |
| $\mathrm{q}=1.0$ | $p=0.1$ | $\rho=0.5$ | $\rho=0.7$ | $p=0$ | $p=0.9$ | $\rho=0.92$ | $p=0.94$ | $p=0.96$ | $\rho=0.98$ |
| $\mathrm{E}\left\{\mathrm{W}_{1}\right\}$ | . 0 | 0.7931 | 1 | 2.6 | 5.5924 | 7.02 | 9.38 | 14.08 | 28.16 |
| $\mathrm{E}\left\{\mathrm{W}_{2}\right\}$ | . 1067 | 1.327 | 3.782 | 7.217 | 18.130 | 23.69 | 32.99 | 51.67 | 107.85 |
| $\mathrm{K}=16$ | $p=0.1$ | $\rho=0.5$ | $\rho=0.7$ | $p=0$ | $p=0.9$ | $\rho=0.92$ | $p=0.94$ | $p=0.96$ | $p=0.98$ |
| $E\left\{W_{1}\right\}$ | . 0 | 0.8001 | 1.8 | 3.3712 | 8.518 | 11.22 | 15.79 | 25.04 | 53.04 |
| $\mathrm{E}\left\{\mathrm{W}_{2}\right\}$ | . 1067 | 1.299 | 3.076 | 4.515 | 6.425 | 6.87 | 7.34 | 7.83 | 8.35 |
| $\mathrm{K}=8$ | $\rho=0.1$ | $p=0.5$ | $\rho=0.7$ | $p=0.8$ | $p=0.9$ | $p=0.92$ | $p=0.94$ | $p=0.96$ | $\rho=0.98$ |
| $E\left\{W_{1}\right\}$ | . 0983 | 0.832 | 2.0468 | 3.7268 | 9.1128 | 11.8 | 16.50 | 25.82 | 53.88 |
| $\mathrm{E}\left\{\mathrm{W}_{2}\right\}$ | . 1067 | 1.1684 | 2.3127 | 3.0928 | 4.0488 | 4.26 | 4.49 | 4.73 | 4.97 |
| $\mathrm{K}=2$ | $p=0.1$ | $\rho=0.5$ | $\rho=0.7$ | $\mathrm{p}=0.8$ | $p=0.9$ | $\rho=0.92$ | $\rho=0.94$ | $p=0.96$ | $p=0.98$ |
| $E\left\{W_{1}\right\}$ | . 0993 | 0.9354 | 2.2931 | 4.0672 | 9.5596 | 12.34 | 16.99 | 26.33 | 54.42 |
| $\mathrm{E}\left\{\mathrm{W}_{2}\right\}$ | . 1030 | 0.7584 | 1.3274 | 1.7311 | 2.2614 | 2.39 | 2.52 | 2.66 | 2.82 |

Table A.4. The model as in table 3: estimations of heavy traffic limits.

| dcp | $\omega_{1}\left(x_{1}\right)$ | $\omega_{2}\left(x_{2}\right)$ | $\omega_{3}\left(x_{3}\right)$ | $\zeta_{0}$ | dcp |  |  | $\omega_{1}\left(x_{1}\right)$ | $\omega_{2}\left(x_{2}\right)$ | $\omega_{3}\left(x_{3}\right)$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 111 | $(3.20)$ | $(9.27)$ | 2.00 | 1.74 | 1.0 | 1.0 | 1.0 | 1.36 | 1.09 | 0.82 |

Table A.5. The model as in table 4: estimations of heavy traffic limits.

| dcp | $\omega_{1}\left(x_{1}\right)$ | $\omega_{2}\left(x_{2}\right)$ | $\omega_{3}\left(x_{3}\right)$ | $\omega_{0}$ | dcp |  |  | $\omega_{1}\left(x_{1}\right)$ | $\omega_{2}\left(x_{2}\right)$ | $\omega_{3}\left(x_{3}\right)$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 111 | 1.01 | 1.13 | 1.25 | 1.13 | 1.0 | 1.0 | 1.0 | 1.59 | 1.27 | 0.95 |

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