

**ELSEVIER**Contents lists available at ScienceDirect

Journal of Combinatorial Theory, Series B

www.elsevier.com/locate/jctb

Infinite primitive and distance transitive directed graphs of finite out-valency [☆]

Daniela Amato ^a, David M. Evans ^b^a *Departamento de Matematica, Universidade de Brasilia, Campus Universitário Darcy Ribeiro, Brasília CEP 70910-900, Brazil*^b *School of Mathematics, University of East Anglia, Norwich NR4 7TJ, UK*

ARTICLE INFO

Article history:

Received 9 August 2012

Available online xxxx

Keywords:

Infinite digraphs
Distance transitivity
High arc transitivity
Primitive groups

ABSTRACT

We give certain properties which are satisfied by the descendant set of a vertex in an infinite, primitive, distance transitive digraph of finite out-valency and provide a strong structure theory for digraphs satisfying these properties. In particular, we show that there are only countably many possibilities for the isomorphism type of such a descendant set, thereby confirming a conjecture of the first Author. As a partial converse, we show that certain related conditions on a countable digraph are sufficient for it to occur as the descendant set of a primitive, distance transitive digraph.

© 2015 Published by Elsevier Inc.

1. Introduction

1.1. Background and main results

We begin with an overview of the paper. Most of the terminology is standard and definitions can be found in the next subsection.

[☆] This work was supported by EPSRC grant EP/G067600/1.

E-mail addresses: d.a.amato@mat.unb.br (D. Amato), d.evans@uea.ac.uk (D.M. Evans).

We are interested in the construction and classification of infinite, vertex transitive directed graphs of finite out-valency whose automorphism groups have additional transitivity properties, such as primitivity, distance transitivity or high arc transitivity. In contrast to the finite case where powerful tools from finite group theory are available, there is no possibility of a complete description of such digraphs. Instead, our results will focus on the structure of the *descendant set* of a vertex in such a digraph: this is the induced subdigraph on the set of vertices reachable from the given vertex by an outward-directed path. Much of the motivation for the work comes from questions of Peter M. Neumann on infinite permutation groups, and work on highly arc transitive digraphs originating in [6].

In [11], Neumann asked whether there exists a primitive permutation group having an infinite suborbit which is paired with a finite suborbit. This amounts to asking whether there is a digraph with infinite in-valency and finite out-valency whose automorphism group is transitive on edges and primitive on vertices. Countable digraphs of this sort were constructed in [8] using amalgamation methods developed in model theory (cf. [5] for background on such methods). In these examples, the descendant sets are directed trees, and the resulting examples are also highly arc transitive. Similar methods were used in [7] to construct continuum-many non-isomorphic countable, primitive, highly arc transitive digraphs all with isomorphic descendant sets. So this suggests that a classification of such digraphs is out of the question, even under the very strong assumption of high arc transitivity. Nevertheless, Neumann (private communication) suggested that a classification of the *descendant sets* in these digraphs might be possible, at least under stronger hypotheses on the automorphism group of the digraph.

Descendant sets in highly arc transitive digraphs of finite out-valency were studied by the first Author in [1,2], following on from results obtained by Möller for locally finite, highly arc transitive digraphs in [10]. This work isolates a small number of quite simple properties (essentially P0, P1, P3 of Section 3 here) satisfied by such descendant sets and shows that these properties have rather strong structural consequences. In particular, the descendant set admits a non-trivial, finite-to-one homomorphism onto a tree. Digraphs having the given properties, but which are not trees are constructed in [2,10]. Moreover (imprimitive) highly arc transitive digraphs having these as descendant sets are constructed in [2–4].

It was conjectured in [2] that there are only countably many directed graphs with the properties for a descendant set isolated in [2]. In Section 2 we reprove some of the results of [2] in a slightly wider context and prove the conjecture. In particular, we have the following (note that a highly arc transitive digraph is distance transitive).

Theorem 1.1. *Suppose that D is a distance transitive digraph of finite out-valency. Assume either that D has infinite in-valency, or that it has no directed cycles. Let $\Gamma = \Gamma_D$ be the descendant set of D . Then there are natural numbers $k(\Gamma)$ and $M(\Gamma)$ with the property that if $\Gamma^{\leq M(\Gamma)}$ denotes the induced subdigraph on the set of vertices in Γ which can*

be reached from the root by a directed path of length at most $M(\Gamma)$, then Γ is determined up to isomorphism by $k(\Gamma)$, $M(\Gamma)$ and the finite digraph $\Gamma^{\leq M(\Gamma)}$.

The proof of this is given at the end of Section 2. Together with Corollary 4.4 of [4], it gives a reasonable picture of the descendant sets in distance transitive digraphs of finite out-valency and infinite in-valency: conditions P0, P1, P3 are necessary and sufficient conditions for a digraph to be a descendant set in such a digraph, and there are only countably many digraphs satisfying these conditions.

In Section 3 we are interested in descendant sets under the additional assumption of primitivity. Then main result is:

Theorem 1.2. *Suppose that a digraph Γ of finite out-valency satisfies conditions P0, P1, P2, P3. Then there is a countable primitive digraph D_Γ of infinite in-valency with descendant set Γ .*

The construction of D_Γ is as in the paper [8], where the descendant set Γ is a tree. However, the proof of primitivity in the general case is much harder than in [8], and this is where the novelty lies in the above result. We do not know whether the condition P2 on Γ is a necessary condition here, but it is satisfied by all of the examples constructed in Section 5 of [2]. So this gives new examples of descendant sets in primitive (and even highly arc transitive) digraphs of finite out-valency and infinite in-valency. It would of course be interesting to find necessary and sufficient conditions on the descendant set in a primitive distance transitive digraph of finite out-valency and infinite in-valency. It would be even more interesting to know whether anything can be said without the assumption of distance transitivity.

1.2. Notation and terminology

A digraph $(D; E(D))$ consists of a set D of vertices, and a set $E(D) \subseteq D \times D$ of ordered pairs of vertices, the (directed) edges. Our digraphs will have no loops and no multiple edges. We will think of a subset X of the set D of vertices as a digraph in its own right by considering the full induced subdigraph on X (so $E(X) = E(D) \cap X^2$). Throughout this paper, ‘subdigraph’ will mean ‘full induced subdigraph’. Thus henceforth, we will not usually distinguish notationally between a digraph and its vertex set. In particular, we will usually refer to the digraph $(D; E(D))$ simply as ‘the digraph D ’. Note that this is a different convention from the usual notation $D = (V(D); E(D))$. Furthermore, we will use notation such as ‘ $\alpha \in D$ ’ to indicate that α is a vertex of the digraph D .

We denote the automorphism group of the digraph D by $\text{Aut}(D)$. We say that D is transitive (respectively, edge transitive) if this is transitive on D (respectively, $E(D)$). We say that D is *primitive* if $\text{Aut}(D)$ is primitive on D , that is, there are no non-trivial $\text{Aut}(D)$ -invariant equivalence relations on D .

The *out-valency* of a vertex $\alpha \in D$ is the size of the set $\{u \in D : (\alpha, u) \in E(D)\}$ of out-vertices of α ; similarly, the *in-valency* of α is the size of the set $\{u \in D : (u, \alpha) \in E(D)\}$ of in-vertices. Let $s \geq 0$ be an integer. An *s-arc* from u to v in D is a sequence $u_0 u_1 \dots u_s$ of $s + 1$ vertices such that $u_0 = u$, $u_s = v$ and $(u_i, u_{i+1}) \in ED$ for $0 \leq i < s$ and $u_{i-1} \neq u_{i+1}$ for $0 < i < s$. Usually our digraphs will be asymmetric, in which case this last condition is redundant. We denote by $D^s(u)$ the set of vertices of D which are reachable by an s -arc from u . The *descendant set* $D(u)$ (or $\text{desc}(u)$) of u is $\bigcup_{s \geq 0} D^s(u)$. Similarly the set $\text{anc}(u)$ of *ancestors* of u is the set of vertices of which u is a descendant.

In particular, fix $\alpha \in D$, and let $\Gamma = D(\alpha)$. If $\text{Aut}(D)$ is transitive on the set of vertices of D , then $D(u) \cong \Gamma$ for all vertices u , and we shall speak of the digraph Γ as *the descendant set* of D .

We say that the digraph D is *highly arc transitive* if for each $s \geq 0$, $\text{Aut}(D)$ is transitive on the set of s -arcs in D . Following [9], we say that a digraph D is (*directed*)-*distance transitive* if for every $s \geq 0$, $\text{Aut}(D)$ is transitive on pairs (u, v) for which there is an s -arc from u to v , but no t -arc for $t < s$. Note that this implies vertex and edge transitivity, but is weaker than being highly arc transitive. We generally exclude the case of null digraphs, where there are no edges.

Henceforth, we shall be interested in the structure of a descendant set $\Gamma = \Gamma(\alpha)$ of a vertex α in some transitive digraph D with finite out-valency m . We will be considering this as a digraph with its full induced structure from D . We refer to α as the *root* of Γ and write $\Gamma = \Gamma(\alpha)$ to indicate that any vertex of Γ is a descendant of α . Similarly, we write Γ^i instead of $\Gamma^i(\alpha)$ for the set of vertices reachable by an i -arc starting at α and if $\beta \in \Gamma(\alpha)$, then we write $\Gamma(\beta) = \text{desc}(\beta) \subseteq \Gamma(\alpha)$. It is clear that if D is highly arc transitive, then $\text{Aut}(\Gamma(\alpha))$ is transitive on s -arcs in $\Gamma(\alpha)$ which start at α . Similarly, if D is distance transitive, then $\text{Aut}(\Gamma(\alpha))$ is transitive on $\Gamma^n(\alpha)$ for each $n \in \mathbb{N}$.

2. The structure of descendant sets

2.1. Preliminaries

We work with digraphs Γ having the following properties:

- G0** $\Gamma = \Gamma(\alpha)$ is a rooted digraph with finite out-valency $m > 0$ and $\Gamma^s(\alpha) \cap \Gamma^t(\alpha) = \emptyset$ whenever $s \neq t$.
- G1** $\Gamma(u) \cong \Gamma$ for all $u \in \Gamma$.
- G2** For $n \in \mathbb{N}$ we have $|\Gamma^n(\alpha)| < |\Gamma^{n+1}(\alpha)|$.
- G3** There is an integer $k \geq 1$ such that if $\ell \geq k$ and $x \in \Gamma^\ell(\alpha)$ and $z \in \Gamma(x)$, then $\text{anc}(z) \cap \Gamma^1(\alpha) = \text{anc}(x) \cap \Gamma^1(\alpha)$.

We shall see that conditions G0, G1, G3 hold when Γ is the descendant set in a distance transitive digraph of finite out-valency and infinite in-valency (Corollary 2.5). The minimum possible k in G3 is the parameter $k(\Gamma)$ which appears in Theorem 1.1. If

$k(\Gamma) = 1$ then Γ is a directed tree, however Section 5 of [2] constructs digraphs $\Gamma(\Sigma, t)$ satisfying G0–G3 with arbitrary value for $k(\Gamma(\Sigma, t))$.

A priori there could be continuum-many isomorphism types of digraphs with these properties. Our main result in this section (Theorem 2.15) is that there are only countably many isomorphism types of digraph Γ which satisfy G0, G1 and G3. To establish this, we show that there is a natural equivalence relation ρ on Γ (refining the ‘layering’ of Γ given by G0) such that the quotient digraph Γ/ρ is a directed tree. If G2 holds then this is not a directed line and the size of the layers Γ^n grows exponentially.

Lemma 2.1. *Suppose D is a (non-null) digraph of finite out-valency which has no directed cycles and is distance transitive. Then any descendant set $\Gamma(\alpha)$ in D satisfies G0.*

Proof. This is the same as the proof of Proposition 3.10 in [6], so we omit the details. \square

Lemma 2.2. *Suppose D is a digraph of finite out-valency with a directed cycle and whose automorphism group is either primitive on vertices or transitive on edges. Then D has finite in-valency.*

Proof. First, suppose that D is edge-transitive. Then there is a K such that every edge of D is in a directed K -cycle. Let $\alpha \in D$. Then every in-vertex β of α is in $D^{K-1}(\alpha)$. But this set is finite, as D has finite out-valency.

Now suppose D is vertex-primitive. Consider the relation \sim on D given by $u \sim v \Leftrightarrow u \in D(v)$ and $v \in D(u)$. This is an $\text{Aut}(D)$ -invariant equivalence relation on D and as D contains a directed cycle, its classes are not singletons. Thus, by primitivity $u \sim v$ for all $u, v \in D$. In particular, every edge of D is contained in a cycle. We can then argue as in the first case. \square

Lemma 2.3. *Suppose Γ satisfies G0, G1 and that for each $i \in \mathbb{N}$ the automorphism group $\text{Aut}(\Gamma)$ is transitive on Γ^i . Then Γ satisfies G3.*

Proof. For $x \in \Gamma^i$, let $t_i = |\text{anc}(x) \cap \Gamma^1|$. By the transitivity assumption, this depends only on i . As $\text{anc}(x) \cap \Gamma^1 \subseteq \text{anc}(z) \cap \Gamma^1$ when $z \in \Gamma(x)$, we have $t_1 \leq t_2 \leq t_3 \leq \dots \leq m$. Choosing k so that t_k is as large as possible, the result follows. \square

Remark 2.4. Note that in the above if G2 also holds, then $t_i < m$. Otherwise, for $\beta \in \Gamma^1$ we have $\Gamma^{i-1}(\beta) = \Gamma^i(\alpha)$ and so $|\Gamma^{i-1}| = |\Gamma^i|$ (by G1), contradicting G2.

Corollary 2.5. *Suppose D is a distance transitive digraph of finite out-valency $m > 0$ and is either of infinite in-valency, or has no directed cycles. Then the descendant set Γ in D satisfies G0, G1, G3. If the automorphism group of D is also primitive on vertices, then $m > 1$ and Γ satisfies G2.*

Proof. By Lemma 2.2, if D has infinite in-valency then D has no directed cycles, so by Lemma 2.1, Γ satisfies G0. As D has transitive automorphism group, G1 holds. Distance transitivity implies that $\text{Aut}(\Gamma)$ is transitive on each Γ^i , so G3 holds.

Suppose $\text{Aut}(D)$ is primitive on vertices of D . If G2 does not hold for some n , then for $\beta, \beta' \in \Gamma^1(\alpha)$ we have $\Gamma^n(\beta) = \Gamma^{n+1}(\alpha) = \Gamma^n(\beta')$. If $m > 1$, then this gives a non-trivial equivalence relation on the vertices of D which is preserved by $\text{Aut}(D)$, and we have a contradiction to primitivity. So it remains to show that $m > 1$. But if $m = 1$, then the underlying (undirected) graph of D has no cycles. This contradicts primitivity of $\text{Aut}(D)$, as it implies that being at even distance in the underlying graph is an equivalence relation on the vertices. \square

2.2. Structure theory

Throughout this section we assume that Γ satisfies G0, G1, G3. We let k be an integer satisfying the condition in G3. The proofs in this section are all adapted from [2].

Lemma 2.6. *Suppose n is a non-negative integer, $\beta \in \Gamma^n(\alpha)$, $\ell \geq k$, $x \in \Gamma^{n+\ell}(\alpha)$ and $z \in \Gamma(x) \cap \Gamma(\beta)$. Then $x \in \Gamma^\ell(\beta)$.*

Proof. This is by induction on n . The case $n = 0$ is trivial as then $\beta = \alpha$. In general let $\gamma \in \Gamma^{n-1}(\alpha)$ be an ancestor of β . By induction hypothesis, $x \in \Gamma^{\ell+1}(\gamma)$. Now work with $\Gamma(\gamma) \cong \Gamma$ (by G1). As $\ell \geq k$ and $z \in \Gamma(x)$ we have $\text{anc}(z) \cap \Gamma^1(\gamma) = \text{anc}(x) \cap \Gamma^1(\gamma)$ (by G3 in $\Gamma(\gamma)$). So $\beta \in \text{anc}(x)$, that is $x \in \Gamma(\beta)$. As $\beta \in \Gamma^n(\alpha)$ and $x \in \Gamma^{n+\ell}(\alpha)$, it follows from G0 that $x \in \Gamma^\ell(\beta)$, as required. \square

Definition 2.7.

(1) Suppose $\beta \in \Gamma$, $x \in \Gamma^n(\beta)$ and $s \leq n$. Define

$$\Gamma_\beta^{-s}(x) = \{w \in \Gamma^{n-s}(\beta) : x \in \Gamma(w)\}.$$

(2) For $\ell \geq k$ and $x, y \in \Gamma^\ell(\alpha)$ write $\rho(x, y)$ iff

$$\Gamma_\alpha^{-k+1}(x) = \Gamma_\alpha^{-k+1}(y).$$

(Say that $\rho(x, y)$ does not hold in all other cases.)

So for $x, y \in \Gamma^\ell(\alpha)$ we have that $\rho(x, y)$ holds iff x, y have the same ancestors in $\Gamma^{\ell-k+1}(\alpha)$. Clearly ρ is an $\text{Aut}(\Gamma)$ -invariant equivalence relation on $\bigcup_{\ell \geq k} \Gamma^\ell$.

Lemma 2.8. *Suppose $\ell \geq k$ and $x, y \in \Gamma^\ell(\alpha)$. If $\Gamma(x) \cap \Gamma(y) \neq \emptyset$, then $\rho(x, y)$ holds.*

Proof. Note that the result holds for $\ell = k$ by G3.

Suppose $\ell = n + k$ with $n \geq 1$ and that $z \in \Gamma(x) \cap \Gamma(y)$. Let $B = \{\beta \in \Gamma^n(\alpha) : z \in \Gamma(\beta)\}$. If $\beta \in B$, then by Lemma 2.6, $x, y \in \Gamma^k(\beta)$. Thus (by the case $\ell = k$ in $\Gamma(\beta)$) we have $\text{anc}(x) \cap \Gamma^1(\beta) = \text{anc}(y) \cap \Gamma^1(\beta)$. But $\Gamma_\alpha^{-k+1}(x), \Gamma_\alpha^{-k+1}(y) \subseteq \bigcup_{\beta \in B} \Gamma^1(\beta)$. Thus $\Gamma_\alpha^{-k+1}(x) = \Gamma_\alpha^{-k+1}(y)$, so $\rho(x, y)$. \square

For $\ell \geq k$ and $x \in \Gamma^\ell(\alpha)$ we write $[x]_\rho$ for the ρ -equivalence class containing x . We use notation such as \mathbf{v}, \mathbf{w} etc. for such classes and write $\Gamma(\mathbf{u}) = \bigcup_{x \in \mathbf{u}} \Gamma(x)$ and $\Gamma^s(\mathbf{u}) = \bigcup_{x \in \mathbf{u}} \Gamma^s(x)$.

Lemma 2.9. Suppose $\ell \geq k$ and $\mathbf{v} \subseteq \Gamma^\ell(\alpha)$ is a ρ -class. Let $w \in \Gamma(\mathbf{v})$. Then $[w]_\rho \subseteq \Gamma(\mathbf{v})$.

Proof. It suffices to prove this when $w \in \Gamma^{\ell+1}(\alpha)$. So suppose that $(v, w), (v', w')$ are directed edges and $\rho(w, w')$ holds. We need to show that $\rho(v, v')$ holds. Let $A = \Gamma_\alpha^{-1}(w)$ and $A' = \Gamma_\alpha^{-1}(w')$. By Lemma 2.8, $A \subseteq [v]_\rho$ and $A' \subseteq [v']_\rho$. By definition, $\Gamma_\alpha^{-k+1}(w) = \bigcup_{a \in A} \Gamma_\alpha^{-k+2}(a)$ and $\Gamma_\alpha^{-k+1}(w') = \bigcup_{a' \in A'} \Gamma_\alpha^{-k+2}(a')$. So $\bigcup_{a \in A} \Gamma_\alpha^{-k+2}(a) = \bigcup_{a' \in A'} \Gamma_\alpha^{-k+2}(a')$, as $\rho(w, w')$ holds. It follows (by taking ancestors one level back) that $\bigcup_{a \in A} \Gamma_\alpha^{-k+1}(a) = \bigcup_{a' \in A'} \Gamma_\alpha^{-k+1}(a')$. But as $A \subseteq [v]_\rho$, the left hand side is equal to $\Gamma_\alpha^{-k+1}(v)$ and similarly the right hand side is equal to $\Gamma_\alpha^{-k+1}(v')$. Thus $\rho(v, v')$ holds. \square

Corollary 2.10. Suppose $\ell \geq k$ and $v \in \Gamma^\ell(\alpha)$. Let \mathbf{v} be the ρ -class containing v . Then the quotient digraph $\Gamma(\mathbf{v})/\rho$ is a rooted directed tree with finite out-valencies.

Proof. The statement follows from Lemmas 2.8 and 2.9. \square

Note that for $\beta \in \Gamma(\alpha)$ we can consider the equivalence relation ρ computed in both $\Gamma(\alpha)$ and $\Gamma(\beta)$, where in the latter we only consider ancestors in $\Gamma(\beta)$ when defining ρ : *a priori* this gives a coarser relation.

Lemma 2.11. Suppose $\beta \in \Gamma^n(\alpha)$ and $x \in \Gamma^\ell(\beta)$ with $\ell \geq 2k - 1$. Then the ρ -class containing x is the same whether it is computed in $\Gamma(\alpha)$ or $\Gamma(\beta)$.

Proof. Note that $x \in \Gamma^{n+\ell}(\alpha)$. First observe that if $y \in [x]_\rho$ (computed in $\Gamma(\alpha)$) then x, y have the same ancestors in $\Gamma^{n+\ell-k+1}(\alpha)$ and so also in $\Gamma^n(\alpha)$: in particular $y \in \Gamma(\beta)$. So to prove the statement, it suffices to show that $\Gamma_\beta^{-k+1}(x) = \Gamma_\alpha^{-k+1}(x)$. It is clear from the definition that $\Gamma_\beta^{-k+1}(x) \subseteq \Gamma_\alpha^{-k+1}(x)$. Conversely, suppose $w \in \Gamma_\alpha^{-k+1}(x)$. Then $w \in \Gamma^{n+\ell-k+1}(\alpha)$ and by assumption $n + \ell - k + 1 \geq n + k$. So by Lemma 2.6 we have $w \in \Gamma(\beta)$ and therefore $w \in \Gamma_\beta^{-k+1}(x)$. \square

Let $\ell \geq 2k - 1$ and let \mathbf{v} be a ρ -class in $\Gamma^\ell(\alpha)$. Let $T(\mathbf{v})$ be the structure consisting of the induced digraph on $\Gamma(\mathbf{v})$ together with the equivalence relation induced by ρ (coming from $\Gamma(\alpha)$). Recall that by Lemma 2.9, $T(\mathbf{v})$ is a union of ρ -classes in $\Gamma(\alpha)$. If \mathbf{w}

is another ρ -class (in $\bigcup_{\ell \geq 2k-1} \Gamma^\ell(\alpha)$) then by a ρ -isomorphism between $T(\mathbf{v})$ and $T(\mathbf{w})$ we mean a digraph isomorphism which respects ρ .

Corollary 2.12. *Suppose \mathbf{v} is a ρ -class in $\Gamma^\ell(\alpha)$ with $\ell \geq 2k - 1$. Then there is a ρ -class \mathbf{w} in $\Gamma^{2k-1}(\alpha)$ and a ρ -isomorphism from $T(\mathbf{w})$ to $T(\mathbf{v})$.*

Proof. Let $v \in \mathbf{v}$ and let $\beta \in \Gamma^{\ell-2k+1}(\alpha)$ be an ancestor of v . So $v \in \Gamma^{2k-1}(\beta)$ and by Lemma 2.11 it follows that $\mathbf{v} \subseteq \Gamma(\beta)$. So $T(\mathbf{v}) \subseteq \Gamma(\beta)$ and the ρ -structure on $T(\mathbf{v})$ is the same whether it is computed in $\Gamma(\alpha)$ or $\Gamma(\beta)$. By G1 there is a digraph isomorphism from $\Gamma(\alpha)$ to $\Gamma(\beta)$, and this induces a ρ -isomorphism between $T(\mathbf{w})$, for some ρ -class $\mathbf{w} \subseteq \Gamma^{2k-1}(\alpha)$, and $T(\mathbf{v}) \subseteq \Gamma^{2k-1}(\beta)$, as required. \square

Thus to any digraph Γ satisfying G0, G1, G3, there are associated a finite number of ρ -isomorphism types of $T(\mathbf{v})$. In particular, we can refine Corollary 2.10 to:

Corollary 2.13. *Suppose $\ell \geq 2k - 1$ and $\mathbf{v} \subseteq \Gamma^\ell(\alpha)$ is a ρ -class. Then the quotient digraph $T(\mathbf{v})/\rho$ is a rooted directed tree with a finite number of out-valencies.* \square

2.3. Counting isomorphism types

We let \mathcal{T} be the class of structures T with the following properties

- T is a digraph of finite out-valency and $T = T(\mathbf{u})$ for some finite set $\mathbf{u} \subseteq T$.
- $T^s(\mathbf{u}) \cap T^t(\mathbf{u}) = \emptyset$ whenever $s \neq t$.
- There is an equivalence relation ρ on T such that each ρ -class is contained in a layer $T^s(\mathbf{u})$.
- The quotient digraph T/ρ is a directed forest.
- For every ρ -class \mathbf{w} there is a ρ -class $\mathbf{v} \subseteq \mathbf{u}$ and a ρ -isomorphism between $T(\mathbf{v})$ and $T(\mathbf{w})$.

We show:

Theorem 2.14. *There are only countably many ρ -isomorphism types of structures in \mathcal{T} .*

Corollary 2.15. *There are only countably many isomorphism types of digraph Γ which have properties G0, G1 and G3.*

Proof. Fix such a Γ . Let T be the disjoint union of digraphs $T(\mathbf{v})$ with the equivalence relation ρ as in the previous section, taking \mathbf{v} to be a ρ -class in Γ^{2k-1} . So in fact, $T = \bigcup_{\ell \geq 2k-1} \Gamma^\ell$. Then $T \in \mathcal{T}$, by Corollaries 2.13 and 2.12. Moreover we can recover Γ from T by looking at the descendant set of any vertex in T . Thus there are only countably many possibilities for Γ , by the above Theorem. \square

We now prove [Theorem 2.14](#). Let $T = T(\mathbf{u}) \in \mathcal{T}$ and let $\mathbf{v}_1, \dots, \mathbf{v}_r$ be the ρ -classes in $T^0 = \mathbf{u}$. We colour a ρ -class \mathbf{v} in T with colour C_i if i is (as small as possible) such that $T(\mathbf{v})$ is ρ -isomorphic to $T(\mathbf{v}_i)$. If $d \in \mathbb{N}$, then we denote by B_T^d the digraph on $\bigcup_{s \leq d} T^s$ together with the structure given by the ρ -classes and the colouring on this set. Similarly if \mathbf{v} is a ρ -class we denote by $B_T^d(\mathbf{v})$ the corresponding structure on $\bigcup_{s \leq d} T^s(\mathbf{v})$.

In the following, by a ρ - C -isomorphism we mean a digraph isomorphism which preserves the relation ρ and the colouring.

Lemma 2.16. *For $T \in \mathcal{T}$ there is a natural number $N = N_T$ with the property that if $d \geq N$, \mathbf{v}, \mathbf{v}' are ρ -classes in T and $\alpha' : B_T^d(\mathbf{v}) \rightarrow B_T^d(\mathbf{v}')$ is a ρ - C -isomorphism, then there is a ρ - C -isomorphism $\alpha : T(\mathbf{v}) \rightarrow T(\mathbf{v}')$ with $\alpha(x) = \alpha'(x)$ for all $x \in \mathbf{v}$.*

Proof. Let A_0 be the group of permutations induced on T^0 by ρ - C -automorphisms of T which fix each ρ -class in T^0 . Similarly for $d \geq 1$ let A_d be the group of permutations induced on T^0 by ρ - C -automorphisms of B_T^d which fix each ρ -class in T^0 . Then $A_d \geq A_{d+1}$ and $A_0 = \bigcap_d A_d$, so there is a smallest integer $N \geq 1$ with $A_N = A_0$. In particular, for any ρ -class \mathbf{v} in T^0 , and $d \geq N$, any permutation of \mathbf{v} which extends to a ρ - C -automorphism of $B_T^d(\mathbf{v})$ extends to an automorphism of $T(\mathbf{v})$. The same is therefore true for any ρ -class in T .

We show that this N has the required property. So let $\mathbf{v}, \mathbf{v}', \dots$ be as in the statement. As \mathbf{v}, \mathbf{v}' have the same colour, there is some ρ - C -isomorphism $\beta : T(\mathbf{v}) \rightarrow T(\mathbf{v}')$. Let β' be its restriction to $B_T^d(\mathbf{v})$. Then α', β' both have image $B_T^d(\mathbf{v}')$ and $\gamma' = (\beta')^{-1} \circ \alpha'$ is a ρ - C -automorphism of $B_T^d(\mathbf{v})$. So as $d \geq N$ there is a ρ - C -automorphism γ of $T(\mathbf{v})$ which agrees with γ' on \mathbf{v} . It is easy to check that $\alpha = \beta \circ \gamma$ is a ρ - C -isomorphism with the required properties. \square

Proposition 2.17. *Suppose $T, S \in \mathcal{T}$ and $d > N_S$. If there is a ρ - C -isomorphism from B_T^d to B_S^d , then there is a ρ - C -isomorphism from B_T^{d+1} to B_S^{d+1} .*

Proof. Let $\Phi : B_T^d \rightarrow B_S^d$ be a ρ - C -isomorphism. Note that $d \geq 1$. Let $\mathbf{v}_1, \dots, \mathbf{v}_s$ be the ρ -classes in T^1 and $\mathbf{w}_i = \Phi(\mathbf{v}_i)$. So $\mathbf{w}_1, \dots, \mathbf{w}_s$ are the ρ -classes in S^1 . For $i \in \{1, \dots, s\}$ there is a ρ -class \mathbf{u}_i in T^0 and a ρ - C -isomorphism $f_i : T(\mathbf{u}_i) \rightarrow T(\mathbf{v}_i)$. Let $\mathbf{z}_i = \Phi(\mathbf{u}_i)$ and $\alpha'_i : B_S^{d-1}(\mathbf{z}_i) \rightarrow B_S^{d-1}(\mathbf{w}_i)$ be given by

$$\alpha'_i(y) = \Phi(f_i(\Phi^{-1}(y))).$$

So α'_i is a ρ - C -isomorphism. As $d-1 \geq N_S$ it follows by [Lemma 2.16](#) that there is a ρ - C -isomorphism $\alpha_i : S(\mathbf{z}_i) \rightarrow S(\mathbf{w}_i)$ which agrees with α'_i on \mathbf{z}_i .

We define $\Psi : B_T^{d+1} \rightarrow B_S^{d+1}$ as follows. For $x \in T^0$ we let $\Psi(x) = \Phi(x)$. If $x \in B_T^{d+1} \setminus T^0$ then there is a unique $i \leq s$ with $x \in B_T^d(\mathbf{v}_i)$ and in this case we define

$$\Psi(x) = \alpha_i(\Phi(f_i^{-1}(x))).$$

It is easy to see that Ψ is a well-defined bijection between B_T^{d+1} and B_S^{d+1} . As f_i, Φ and α_i all preserve ρ -classes and the colouring, the same is true of Ψ . So it remains to show that Ψ preserves edges and non-edges.

First we show that if $x \in B_T^1$, then $\Psi(x) = \Phi(x)$. If $x \in T^0$ then this is by definition of Ψ . If $x \in T^1$ then $x \in \mathbf{v}_i$ for some unique $i \leq s$. So $f_i^{-1}(x) \in \mathbf{u}_i$ and $\Phi(f_i^{-1}(x)) \in \mathbf{z}_i$, whence

$$\Psi(x) = \alpha_i \Phi f_i^{-1}(x) = \alpha'_i \Phi f_i^{-1}(x) = \Phi(x).$$

Thus Ψ preserves edges and non-edges in B_T^1 .

If $x, y \in B_T^{d+1} \setminus T^0$ and (x, y) is an edge, then $x, y \in B_T^d(\mathbf{v}_i)$ for some i . Then $\Psi(x) = \alpha_i \Phi f_i^{-1}(x)$ and $\Psi(y) = \alpha_i \Phi f_i^{-1}(y)$ and so, as α_i, Φ and f_i preserve edges, $(\Psi(x), \Psi(y))$ is an edge in B_S^d . By the same argument, if $x, y \in B_T^d(\mathbf{v}_i)$ and (x, y) is a non-edge, then $(\Psi(x), \Psi(y))$ is a non-edge. Finally, if x, y lie in different $B_T^d(\mathbf{v}_i)$ then $\Psi(x), \Psi(y)$ lie in different $B_S^d(\mathbf{w}_i)$, so $(\Psi(x), \Psi(y))$ is a non-edge. \square

Corollary 2.18. *Suppose $T, S \in \mathcal{T}$ and B_T^d, B_S^d are ρ - C -isomorphic for some $d \geq N_S$. Then T and S are ρ - C -isomorphic.*

Proof. By assumption and [Proposition 2.17](#), for $n \geq N_S$ the set I_n of ρ - C -isomorphisms $B_T^n \rightarrow B_S^n$ is non-empty. Restriction gives a map $I_{n+1} \rightarrow I_n$ and so, as each I_n is finite, König's Lemma implies that there is a ρ - C -isomorphism $T \rightarrow S$. \square

Proof of Theorem 2.14. Suppose $T \in \mathcal{T}$. As above, consider this with a colouring of the ρ -classes. Let $N = N_T$ be as in [Lemma 2.16](#). Then by [Corollary 2.18](#), the (coloured) ball B_T^{N+1} determines T within \mathcal{T} up to isomorphism. There are only countably many possibilities for this finite structure, hence the result. \square

We now give the proof of [Theorem 1.1](#).

Proof of Theorem 1.1. Let D and $\Gamma = \Gamma_D$ be as in the statement of the theorem. First, we define the numbers $k(\Gamma)$ and $M(\Gamma)$. By [Corollary 2.5](#), Γ satisfies G0, G1, G3. Let $k = k(\Gamma)$ be the smallest value of k which satisfies G3 for Γ .

Let \mathbf{w} be a ρ -class in Γ^{2k-1} . By distance transitivity, $\text{Aut}(\Gamma)$ is transitive on each Γ^ℓ , so if $\ell \geq 2k-1$ and \mathbf{v} is a ρ -class in Γ^ℓ , then there is a ρ -isomorphism from $\Gamma(\mathbf{w})$ to $\Gamma(\mathbf{v})$, by [Corollary 2.12](#). Let $T_\Gamma \in \mathcal{T}$ consist of $\Gamma(\mathbf{w})$ together with its ρ -structure. Note that we have only one 'colour' C_i used here, so ρ -isomorphisms will be ρ - C -isomorphisms in what follows.

Let $M(\Gamma) = 2k(\Gamma) + N_{T_\Gamma}$, where N_{T_Γ} is as in [Lemma 2.16](#) (chosen as small as possible). Thus, from the proof of [Lemma 2.16](#), N_{T_Γ} is the smallest value of N such that any permutation of \mathbf{w} which extends to a ρ -automorphism of $T_\Gamma^{\leq N}$ extends to a ρ -automorphism of T_Γ .

Now suppose that D_1, D_2 are distance transitive digraphs satisfying the hypotheses of the theorem. Let $\Gamma_i = \Gamma_{D_i}$ and suppose that $k(\Gamma_1) = k(\Gamma_2) = k$, $M(\Gamma_1) = M(\Gamma_2) = M$ and $\theta : \Gamma_1^{\leq M} \rightarrow \Gamma_2^{\leq M}$ is an isomorphism. As $k(\Gamma_1) = k(\Gamma_2)$, θ gives a ρ -isomorphism $\Gamma_1^\ell \rightarrow \Gamma_2^\ell$ for $k \leq \ell \leq M$.

Let $\mathbf{w}_i \in \Gamma_i^{2k-1}$ be ρ -classes, with $\theta(\mathbf{w}_1) = \mathbf{w}_2$. Let $T_i = \Gamma_i(\mathbf{w}_i)$, considered also with its ρ -structure. Then $N_{T_1} = N_{T_2} = N$ and θ gives a ρ -isomorphism between the balls $B_{T_1}^{N+1}$ and $B_{T_2}^{N+1}$. By [Corollary 2.18](#) (and the above remark on colours) T_1 and T_2 are isomorphic. It then follows that Γ_1 and Γ_2 are isomorphic, as required. \square

3. Constructions

In this section, we prove [Theorem 1.2](#). The construction of the digraphs D_Γ is as in [\[8\]](#) and we recall briefly some notation and terminology from there.

Suppose D is a digraph and $A \subseteq D$. We write $A \leq D$ if for every $a \in A$ we have $\text{desc}(a) \subseteq A$. We say that $A \leq D$ is *finitely generated* (f.g.) if there is a finite $X \subseteq A$ with $A = \bigcup_{a \in X} \text{desc}(a)$, and say that X is a *generating set* for A .

We write $A \leq^+ D$ if $A \leq D$ and

- (i) for every $b \in D$, if $\text{desc}(b) \setminus A$ is finite, then $b \in A$;
- (ii) for all $b \in D$, $\text{desc}(b) \cap A$ is finitely generated.

It is easy to check (cf. Lemma 2.2 of [\[8\]](#)) that if $A \leq^+ B \leq^+ C$ then $A \leq^+ C$ (and similarly for \leq).

We work with a digraph Γ having the following properties:

- P0** $\Gamma = \Gamma(\gamma)$ is a rooted digraph with finite out-valency $m > 0$ and $\Gamma^s(\gamma) \cap \Gamma^t(\gamma) = \emptyset$ whenever $s \neq t$.
- P1** $\Gamma(u) \cong \Gamma$ for all $u \in \Gamma$.
- P2** For all $a \in \Gamma$ we have $\text{desc}(a) \leq^+ \Gamma$.
- P3** For all natural numbers n , $\text{Aut}(\Gamma)$ is transitive on Γ^n .

Of course, P0 and P1 are the same as G0, G1 and P3 implies G3 (as in [Lemma 2.3](#)). From Section 2.1, if Γ is the descendant set of a vertex in an infinite, distance transitive digraph D of finite out-valency and with no directed cycles then Γ satisfies P0, P1, P3. If D is primitive, then (as noted in [\[2\]](#) under the stronger assumption of high arc transitivity) Γ satisfies the following weaker version of P2:

- P2'** For all $a_1, a_2 \in \Gamma$, if $\Gamma(a_1) \setminus \Gamma(a_2)$ and $\Gamma(a_2) \setminus \Gamma(a_1)$ are finite, then $a_1 = a_2$.

Note in particular that P2 implies that different vertices have different sets of out-vertices.

Section 5 of [2] gives examples $\Gamma(\Sigma, k)$ which satisfy P0, P1, P2', P3 and it can be checked that these examples also satisfy P2. In this section we prove that if Γ satisfies P0–P3, then there is a primitive digraph D_Γ with Γ as its descendant set. If Γ has the property that $\text{Aut}(\Gamma)$ is transitive on n -arcs from γ (as is the case with the $\Gamma(\Sigma, k)$ from [2]), then the D_Γ which we construct will be highly arc transitive.

The construction of D_Γ is essentially the same Fraïssé amalgamation class construction which was used in [8]. We will recall this briefly, making use of results from [4]. Once we have D_Γ , the main work of the section will be in proving primitivity of $\text{Aut}(D_\Gamma)$.

So suppose Γ satisfies P0–P3. Let $\bar{\mathcal{C}}_\Gamma$ consist of the digraphs A with the property that for every $a \in A$, $\text{desc}(a) \leq^+ A$ and $\text{desc}(a) \cong \Gamma$. Let \mathcal{C}_Γ be the finitely generated elements of $\bar{\mathcal{C}}_\Gamma$. Note that $\Gamma \in \mathcal{C}_\Gamma$, so in particular, \mathcal{C}_Γ is non-empty.

If $A, B \in \bar{\mathcal{C}}_\Gamma$ a digraph embedding $f : A \rightarrow B$ is called a \leq^+ -embedding if $f(A) \leq^+ B$. We say that \leq^+ -embeddings $f_i : A \rightarrow B_i$ (for $i = 1, 2$) are *isomorphic* if there is a digraph isomorphism $h : B_1 \rightarrow B_2$ with $f_2 = h \circ f_1$.

Lemma 3.1. (Cf. 2.14 of [8].) *Suppose Γ satisfies P0–P3. Then*

- (1) *there are countably many isomorphism types of digraphs in \mathcal{C}_Γ ;*
- (2) *if $A, B \in \mathcal{C}_\Gamma$ then there are countably many isomorphism types of \leq^+ -embeddings $f : A \rightarrow B$.*

Proof. This follows from results in Section 4 of [4]. The digraph Γ satisfies the conditions T1, T2, T3, T4 in Theorem 4.3 of [4] (the first three are just P0, P1, P3 and T4 follows from these as in Remark 4 of [4]). As in the proof of Corollary 4.4 of [4], it follows that Γ satisfies conditions (C1), (C2) of Theorem 4.1 of [4]. Under these conditions, Lemma 4.2 of [4] gives the stronger result that there are countably many isomorphism types of digraph embeddings $f : A \rightarrow B$ with $A, B \in \mathcal{C}_\Gamma$ and $f(A) \leq B$. (Note that \mathcal{C}_Γ as defined here is a subset of the \mathcal{C}_Γ defined in Section 4.1 of [4].) The result we want follows: for (1), take $A = \emptyset$ and (2) is immediate. \square

It is easy to show that \mathcal{C}_Γ is closed under free amalgamation over finitely generated \leq^+ -subsets (as in Lemma 2.6 of [8]). More formally, if $B_1, B_2 \in \mathcal{C}_\Gamma$ and $A \leq^+ B_i$ is f.g., then the digraph F which has vertices the disjoint union of B_1 and B_2 over A and whose edges are the edges of B_1 and B_2 is also in \mathcal{C}_Γ . Furthermore, $B_1, B_2 \leq^+ F$. This gives the following \leq^+ -amalgamation property:

Lemma 3.2. *Suppose $A, B_1, B_2 \in \mathcal{C}_\Gamma$ and $f_i : A \rightarrow B_i$ are \leq^+ -embeddings (for $i = 1, 2$). Then there exist $F \in \mathcal{C}_\Gamma$ and \leq^+ -embeddings $g_i : B_i \rightarrow F$ with the property that $g_1(f_1(a)) = g_2(f_2(a))$ for all $a \in A$. \square*

Note that if f_1 is inclusion, then we can also take g_1 to be inclusion here.

Once we have these lemmas, the following existence and uniqueness result is fairly standard and we omit some of the details of the proof.

Theorem 3.3. *There is a countable digraph D_Γ with the properties:*

- (1) *If $a \in D_\Gamma$ then $\text{desc}(a) \leq^+ D_\Gamma$ and $\text{desc}(a) \cong \Gamma$.*
- (2) *If $X \subseteq D_\Gamma$ is finite, there is a f.g. $A \leq^+ D_\Gamma$ with $X \subseteq A \in \mathcal{C}_\Gamma$.*
- (3) *If $A \leq^+ D_\Gamma$ is f.g. and $f : A \rightarrow B \in \mathcal{C}_\Gamma$ is such that $f(A) \leq^+ B$ then there is $g : B \rightarrow D_\Gamma$ with $gf(a) = a$ for all $a \in A$ and $g(B) \leq^+ D_\Gamma$.*

Moreover, D_Γ is uniquely determined up to isomorphism by these conditions and is \leq^+ -homogeneous, meaning that if $A_1, A_2 \leq^+ D_\Gamma$ are f.g. and $h : A_1 \rightarrow A_2$ is an isomorphism, then h extends to an automorphism of D_Γ .

Proof. Note that (1) here follows from (2). For the existence part, we build a chain of digraphs $D_i \in \mathcal{C}_\Gamma$

$$D_1 \leq^+ D_2 \leq^+ D_3 \leq^+ \dots$$

with the property:

- (*) if $A \leq^+ D_i$ is finitely generated and $f : A \rightarrow B$ is a \leq^+ -embedding with $B \in \mathcal{C}_\Gamma$, then there is some $j \geq i$ and a \leq^+ -embedding $g : B \rightarrow D_j$ such that $g(f(a)) = a$ for all $a \in A$.

Once we have this, we let D_Γ be the union $\bigcup_{n \in \mathbb{N}} D_n$. Then (2) follows as each D_n is in \mathcal{C}_Γ , and (3) follows from (*).

In order to obtain (*) we build the D_n inductively. During this process, there will be countably many ‘tasks’ to be performed: there are countably many choices of f.g. A in each D_i and countably many isomorphism types of \leq^+ -embeddings $f : A \rightarrow B$ with $B \in \mathcal{C}_\Gamma$ (by Lemma 3.1). As we have countably many steps available during the construction, it will suffice to show how to complete one of these tasks: ensuring that they are all completed during some stage of the construction is then just a matter of organisation (see the proof of Theorem 2.8 of [7] for a formal way of doing this).

So suppose D_n has been constructed, $A \leq^+ D_n$ is f.g. and $f : A \rightarrow B$ is a \leq^+ -embedding with $B \in \mathcal{C}_\Gamma$. Using amalgamation (Lemma 3.2) we can find $D_n \leq^+ D_{n+1}$ and a \leq^+ -embedding $g : B \rightarrow D_{n+1}$ with $g(f(a)) = a$ for all $a \in A$, as required.

This completes the construction of some countable digraph D_Γ with properties (1), (2), (3). For the ‘Moreover’ part, suppose D'_Γ is also a countable digraph with properties (1), (2), (3). Suppose $A \leq^+ D_\Gamma$ and $A' \leq^+ D'_\Gamma$ are f.g. and $h : A \rightarrow A'$ is an isomorphism. It will suffice to prove that h extends to an isomorphism $D_\Gamma \rightarrow D'_\Gamma$. As D_Γ, D'_Γ are countable, this follows by a back-and-forth argument (and symmetry) once we show:

Claim. *If $c \in D_\Gamma$ there exist finitely generated $B \leq^+ D_\Gamma$ and $B' \leq^+ D'_\Gamma$ with $A \leq^+ B$, $A' \leq^+ B'$ and $c \in B$ and an isomorphism $g : B \rightarrow B'$ extending h .*

Existence of B here follows from (2) in D_Γ . Existence of g and B' then follows from (3) in D'_Γ (applied to $h^{-1} : A' \rightarrow B$). \square

It is clear that \leq^+ -homogeneity together with property (1) in [Theorem 3.3](#) imply that $\text{Aut}(D_\Gamma)$ is transitive on vertices. Moreover, for every $a \in D_\Gamma$, any automorphism of the descendent set $D_\Gamma(a)$ extends to an automorphism of D_Γ (necessarily fixing a) and so by P3, $\text{Aut}(D_\Gamma/a)$ (the stabiliser of a) is transitive on $D_\Gamma^n(a)$ (vertices reachable by an n -arc from a). Thus D_Γ is distance transitive.

The remainder of this section is devoted to showing:

Theorem 3.4. *With the above notation, $\text{Aut}(D_\Gamma)$ is primitive on the vertices of D_Γ .*

By the above remarks, [Theorem 1.2](#) then follows.

The following lemma is a simple application of free amalgamation and the extension property (3) in [Theorem 3.3](#), but we shall give the details.

Lemma 3.5. *Suppose $A \leq^+ B \leq^+ D_\Gamma$ and A, B are finitely generated. Suppose h is an automorphism of A . Then h can be extended to $g \in \text{Aut}(D_\Gamma)$ so that $B \cap gB = A$.*

Proof. Let B' be any set with $B \cap B' = A$ and $|B \setminus A| = |B' \setminus A|$. Extend h to a bijection $s : B \rightarrow B'$. Let E be the set of directed edges in B and define a digraph relation E' on B' by $(x, y) \in E' \Leftrightarrow (s^{-1}x, s^{-1}y) \in E$. As s restricted to A is h and this is a digraph isomorphism, we have $E \cap A^2 = E' \cap A^2$. Clearly s is then a digraph isomorphism. Let $F = B \cup B'$ with digraph relation $E \cup E'$. So F is the free amalgam of B and B' over A and $B, B' \leq^+ F \in \mathcal{C}_\Gamma$. By the extension property ([Theorem 3.3\(3\)](#)) over B , we can assume that $F \leq^+ D_\Gamma$. Then $s : B \rightarrow B'$ extends to an automorphism g of D_Γ (by the ‘Moreover’ in [Theorem 3.3](#)), and this has the required properties. \square

If $a, b \in D_\Gamma$ and $a \neq b$, let $\Delta(a, b)$ be the orbital digraph which has (a, b) as an edge. So this is the digraph with vertex set D_Γ and edge set the $\text{Aut}(D_\Gamma)$ -orbit which contains (a, b) . By D.G. Higman’s criterion, to prove the primitivity, it will be enough to show that each such $\Delta(a, b)$ is connected (meaning that its underlying undirected graph is connected).

Lemma 3.6. *If (a, b) and (a, b') are in the same $\text{Aut}(D_\Gamma)$ -orbit, and $\Delta(b, b')$ is connected, then $\Delta(a, b)$ is connected (of diameter at most twice that of $\Delta(b, b')$).*

Proof. There is an (undirected) $\Delta(a, b)$ -path bab' from b to b' . Thus if (b_1, b_2) is an edge in $\Delta(b, b')$ there is a $\Delta(a, b)$ -path of length 2 from b_1 to b_2 . As $\Delta(b, b')$ is connected, if $x, y \in D_\Gamma$ there is a $\Delta(b, b')$ -path from x to y , and therefore there is a $\Delta(a, b)$ -path from x to y with at most twice as many edges. \square

Lemma 3.7. *If $a, b \in D_\Gamma$ are distinct vertices and $\text{desc}(a) \cap \text{desc}(b) = \emptyset$, then $\Delta(a, b)$ is connected (of diameter at most 4).*

Proof. First suppose that $\text{desc}(a) \cup \text{desc}(b) \leq^+ D_\Gamma$. Given $x_1, x_2 \in D_\Gamma$, by [Theorem 3.3\(2\)](#) there is a finitely generated $Z \leq^+ D_\Gamma$ with $x_1, x_2 \in Z$. Let X be the disjoint union of Z and a copy C of Γ . So X is the free amalgam of Z and C over the empty set and $X \in \mathcal{C}_\Gamma$. By [Theorem 3.3\(3\)](#) we may assume that $X \leq^+ D_\Gamma$. Let $c \in D_\Gamma$ be the generator of C . Then $\text{desc}(c) \cap \text{desc}(x_i) = \emptyset$ and $\text{desc}(x_i) \cup \text{desc}(c) \leq^+ D_\Gamma$ (for $i = 1, 2$). It follows (using the ‘Moreover’ part of [Theorem 3.3](#)) that (c, x_1) and (c, x_2) are edges in $\Delta(a, b)$. So $\Delta(a, b)$ has diameter 2.

In general, there is $b' \in D_\Gamma$ such that (a, b') is an edge in $\Delta(a, b)$ and f.g. $B, B' \leq^+ D_\Gamma$ with $a, b \in B$ and $a, b' \in B'$ satisfying $B \cup B' \leq^+ D_\Gamma$ and $B \cap B' = \text{desc}(a)$ (again, this uses free amalgamation and property (3) of D_Γ ; alternatively we can apply [Lemma 3.5](#)). Then b, b' are as in the first part of the proof and so $\Delta(b, b')$ has diameter 2. Therefore by [Lemma 3.6](#), $\Delta(a, b)$ has diameter at most 4. \square

The main work is in proving:

Proposition 3.8. *Suppose $a, b \in D_\Gamma$ are distinct vertices and $a \notin \text{desc}(b)$ and $b \notin \text{desc}(a)$. Then there exist $r \in \mathbb{N}$ and $g_1, \dots, g_r \in \text{Aut}(D_\Gamma)$ such that, setting $b_0 = a$, $b_1 = b$ and $b_{i+1} = g_i b_i$ for $i \geq 1$, we have $g_i \in \text{Aut}(D_\Gamma/b_{i-1})$ and $\text{desc}(b_r) \cap \text{desc}(b_{r+1}) = \emptyset$.*

We postpone the proof for now, and carry on with the proof of [Theorem 3.4](#).

Lemma 3.9. *Suppose $a, b \in D_\Gamma$ are distinct vertices and $a \notin \text{desc}(b)$ and $b \notin \text{desc}(a)$. Then $\Delta(a, b)$ is connected.*

Proof. Let r, g_i and b_i be as in [Proposition 3.8](#). By [Lemma 3.7](#), $\Delta(b_r, b_{r+1})$ is connected. So by [Lemma 3.6](#), $\Delta(b_{r-1}, b_r)$ is connected. Proceeding in this way, we obtain that $\Delta(b_0, b_1)$ is connected. \square

The remaining case to consider is:

Lemma 3.10. *Suppose $a \in \text{desc}(b)$ (and $a \neq b$). Then $\Delta(a, b)$ is connected.*

Proof. By [Lemma 3.5](#), there exists $g \in \text{Aut}(D_\Gamma/a)$ with $\text{desc}(b) \cap \text{desc}(gb) = \text{desc}(a)$. By [Lemma 3.9](#), $\Delta(b, gb)$ is connected. So by [Lemma 3.6](#), $\Delta(a, b)$ is connected. \square

By D.G. Higman’s criterion, [Lemmas 3.9 and 3.10](#) establish [Theorem 3.4](#). Thus, it remains to prove [Proposition 3.8](#). We first recall some definitions and results from [Section 2](#).

We tend to write $\Gamma(a)$ for $\text{desc}(a)$ in order to emphasise the isomorphism with Γ . We write $\Gamma^n(a)$ for the descendants reachable by an n -arc, and say that $b \in \Gamma^n(a)$ is *at level n* with respect to a . Variations such as $\Gamma^{\geq n}(a)$ (for $\bigcup_{m \geq n} \Gamma^m(a)$) will also be used.

For any $x \in \Gamma^\ell$ (with $\ell \geq 1$) consider the number of ancestors of x in Γ^1 . By P3 this depends only on ℓ and is non-decreasing as ℓ increases. Thus there is some value k of ℓ for which it reaches a maximum size (which is necessarily less than q , the out-valency of Γ). Let $K = 2k - 1$.

Now we work in D_Γ . For $a, x, y \in D_\Gamma$ with $x, y \in \Gamma^\ell(a)$ and $\ell \geq k$, we write $\rho_a(x, y)$ to indicate that x, y have the same ancestors in $\Gamma^{\ell-k+1}(a)$. So this is the relation ρ on $\Gamma(a)$ used in Section 2. This is an equivalence relation on $\Gamma^\ell(a)$ which is clearly invariant under the stabiliser $\text{Aut}(D_\Gamma/a)$. Denote the equivalence class containing x by $\rho_a[x]$.

Lemma 3.11.

- (1) If $x \in \Gamma(a)$ and $y \in \Gamma^{\geq K}(x)$ and $\rho_a(z, y)$, then $z \in \Gamma(x)$.
- (2) If $x \in \Gamma^m(a)$ and $\ell \geq K$, then $\Gamma(x) \cap \Gamma^{m+\ell}(a)$ is a union of ρ_a -classes.
- (3) If $x, y \in \Gamma^\ell(a)$ and $\ell \geq k$ and $\Gamma(x) \cap \Gamma(y) \neq \emptyset$ then $\rho_a(x, y)$.

Proof. (1) This follows directly from Lemma 2.11.

(2) This follows immediately from (1).

(3) This is Lemma 2.8. \square

Let $\ell \geq k$ and $a \in D_\Gamma$. For $x, y \in \Gamma^\ell(a)$ write $\sigma_a^0(x, y)$ if $\Gamma(x) \cap \Gamma(y) \neq \emptyset$ and let σ_a be the transitive closure of this relation. Note that this is an equivalence relation (on each $\Gamma^\ell(a)$) which is clearly $\text{Aut}(D_\Gamma/a)$ -invariant. Moreover, by (3) above, $\sigma_a(x, z)$ implies $\rho_a(x, z)$. Write $\sigma_a[x]$ for the σ_a -class containing x .

Remark 3.12. In the examples in [2], we have $\rho_a = \sigma_a$. It is not clear whether this holds for arbitrary Γ satisfying P0–P3.

Lemma 3.13. Suppose $b, b' \in D_\Gamma$ are such that if $x \in X = \Gamma(b) \cap \Gamma(b')$ then $x \in \Gamma^\ell(b) \Leftrightarrow x \in \Gamma^\ell(b')$ (so points in the intersection are at the same ‘level’ with respect to b and b'). Then there exists $n \geq K$ such that $X = \text{desc}(X \cap \Gamma^{\leq n-K}(b))$. Moreover, for such an n , we have $X \cap \Gamma^n(b) = X \cap \Gamma^n(b')$ and if $y \in X \cap \Gamma^n(b)$, then $\sigma_b[y] = \sigma_{b'}[y]$.

Proof. As X is finitely generated, it will suffice to take n large enough so that $\Gamma^{\leq n-K}$ contains a generating set for X . For the rest, note first that if $y \in X \cap \Gamma^n(b)$, then by the opening assumption on levels, $y \in X \cap \Gamma^n(b')$. Moreover, the assumption on n means that there is $x \in X$ such that $y \in \Gamma^{\geq K}(x)$. So by Lemma 3.11(2), $\rho_b[y] \subseteq X$. Thus $\sigma_b[y] \subseteq X$. Similarly $\sigma_{b'}[y] \subseteq X$. Now, if $\sigma_b^0(y, z)$ holds then $z \in X$ and so (by definition of σ^0) also $\sigma_{b'}^0(y, z)$ holds. The statement follows. \square

Proof of Proposition 3.8. Let $X = \Gamma(a) \cap \Gamma(b)$. Clearly we may assume that X is non-empty. Then $X \leq^+ \Gamma(a)$ (as $\Gamma(b) \leq^+ D_\Gamma$) and $X \neq \Gamma(a)$. By free amalgamation (over $\Gamma(a)$) there is a copy b' of b over $\Gamma(a)$ with $\Gamma(b) \cap \Gamma(b') = X$. More precisely, apply Lemma 3.5 with $A = \Gamma(a)$, h the identity on A , and $B \leq^+ D_\Gamma$

finitely generated with $a, b \in B$. It follows that there is g_1 in the pointwise stabiliser $\text{Aut}(D_\Gamma/\Gamma(a))$ of $\Gamma(a)$ such that $g_1(B) \cap B = \Gamma(a)$. Then, with $b' = g_1b$, we have $X \subseteq \Gamma(b) \cap \Gamma(b') \subseteq \Gamma(b) \cap B \cap g_1(B) = \Gamma(b) \cap \Gamma(a) = X$ as g_1 fixes all elements of X . Moreover, because g_1 fixes all elements of X and sends b to b' , the points in X are at the same level with respect to b and b' , as in [Lemma 3.13](#).

Let $b_0 = a$, $b_1 = b$ and $b_2 = b'$. Suppose, for $i \geq 2$, that $g_i \in \text{Aut}(D_\Gamma/b_{i-1})$ and $b_{i+1} = g_i b_i$ is such that $\Gamma(b_{i+1}) \cap \Gamma(b_i) \subseteq \Gamma(b_{i-1})$ and $b_{i+1} \neq b_i$ (the existence of such g_i will be proved later on). We make a series of claims, refining the choice of the g_i as we proceed. We may suppose $\Gamma(b_1) \cap \Gamma(b_2) \neq \emptyset$.

Claim 1. *For $i \geq 1$, if $x \in \Gamma(b_i) \cap \Gamma(b_{i+1})$, then x is at the same level with respect to b_i and b_{i+1} .*

Proof. We prove this by induction, noting first that it holds for $i = 1$, by construction. By assumption on the g_i , we have $x \in \Gamma(b_{i-1}) \cap \Gamma(b_i) \cap \Gamma(b_{i+1})$. So by inductive assumption, x has the same level with respect to b_{i-1} and b_i . But $\Gamma(b_{i-1}) \cap \Gamma(b_{i+1})$ is $g_i(\Gamma(b_{i-1}) \cap \Gamma(b_i))$, so a point in this intersection has the same level with respect to b_{i-1} and b_{i+1} . Thus x has the same level with respect to b_{i-1}, b_i and b_{i+1} . \square

Choose $n \geq K$ such that $Y_1 = \Gamma(b_1) \cap \Gamma(b_2)$ is generated by its intersection with $\Gamma^{\leq n-K}(b_1)$. Let $Z_1 = Y_1 \cap \Gamma^n(b_1) = \Gamma^n(b_1) \cap \Gamma^n(b_2)$. By [Lemma 3.13](#) we have:

Claim 2. *Z_1 is a union of sets which are simultaneously σ_{b_1} -classes and σ_{b_2} -classes.*

Claim 3. *For $i \geq 2$, if $Z_i = \Gamma^n(b_i) \cap \Gamma^n(b_{i+1})$, then Z_i is a union of sets which are simultaneously σ_{b_j} -classes for all $1 \leq j \leq i+1$.*

Proof. Note that [Claim 1](#) and the fact that $\Gamma(b_{i+1}) \cap \Gamma(b_i) \subseteq \Gamma(b_{i-1})$ imply that $Z_i \subseteq Z_{i-1} \subseteq \dots \subseteq Z_1$. By Claim 2, we can assume inductively that Z_{i-1} is a union of subsets which are simultaneously σ_{b_j} -classes for $1 \leq j \leq i$. In particular, Z_{i-1} is a union of sets which are simultaneously $\sigma_{b_{i-1}}$ and σ_{b_i} -classes. As $g_i \in \text{Aut}(D_\Gamma/b_{i-1})$ and $g_i b_i = b_{i+1}$ it follows that $g_i Z_{i-1}$ is a union of sets which are simultaneously $\sigma_{b_{i-1}}$ and $\sigma_{b_{i+1}}$ -classes. Thus $Z_i = Z_{i-1} \cap g_i Z_{i-1}$ is a union of $\sigma_{b_{i-1}}$ -classes, and all of these classes are also σ_{b_j} -classes for $1 \leq j \leq i$ and $j = i+1$, as required. \square

Claim 4. *For $i \geq 1$ we have that $Y_i = \Gamma(b_i) \cap \Gamma(b_{i+1})$ is generated by its intersection with $\Gamma^{\leq n}(b_i)$.*

Proof. For $i = 1$ this is by choice of n . So suppose $i \geq 2$. Then we have $Y_i = Y_{i-1} \cap g_i Y_{i-1}$ and $Y_{i-1}, g_i Y_{i-1} \subseteq \Gamma(b_{i-1})$. Suppose $y \in Y_i$ is at level $m > n$ in $\Gamma(b_{i-1})$ (and therefore also in $\Gamma^m(b_i)$ and $\Gamma^m(b_{i+1})$, by [Claim 1](#)). Then there exist $x \in Y_{i-1}$ and $x' \in g_i Y_{i-1}$ which are at level n (in $\Gamma(b_{i-1})$) and which are ancestors of y . Then, because of y , we have

that $\sigma_{b_{i-1}}(x, x')$ holds. But $x \in Z_{i-1}$ and Z_{i-1} is a union of $\sigma_{b_{i-1}}$ -classes, so $x' \in Z_{i-1}$ and therefore $x' \in Y_i$. So y has an ancestor in Z_i , as required. \square

Suppose now that we can choose the g_i so that if $Z_i \neq \emptyset$ then $g_i Z_i \neq Z_i$, and therefore Z_{i+1} is a proper subset of Z_i . As these sets are finite, for some r we must have $Z_r = \emptyset$. It then follows from [Claim 4](#) that $Y_r = \emptyset$, as required by the proposition.

It remains to explain how to construct the g_i , for $i \geq 2$. We do this inductively, so suppose we have constructed up to g_{i-1} and have b_1, \dots, b_i with the required properties. Suppose $Z_{i-1} \neq \emptyset$. We need to find $g_i \in \text{Aut}(D_\Gamma/b_{i-1})$ such that (with the above notation) $Y_i \subseteq \Gamma(b_{i-1})$ and $g_i Z_{i-1} \neq Z_{i-1}$. First, note that Z_{i-1} is a proper subset of $\Gamma^n(b_{i-1})$. If not, then $\Gamma^n(b_{i-1}) \subseteq \Gamma^n(b_i)$ and as these are finite sets of the same size, we obtain that $\Gamma^n(b_{i-1}) = \Gamma^n(b_i)$. So b_{i-1} and b_i have the same descendants in level n and as $b_{i-1} \neq b_i$, this contradicts property (1) in the definition of D_Γ in [Theorem 3.3](#). It then follows by property P3 of Γ that there is an automorphism h of $\Gamma(b_{i-1})$ with $h Z_{i-1} \neq Z_{i-1}$. We claim that h extends to an automorphism g_i of D_Γ with the property that $\Gamma(g_i b_i) \cap \Gamma(b_i) \subseteq \Gamma(b_{i-1})$, and this will suffice.

To do this, let $B \leq^+ D_\Gamma$ be finitely generated and contain b_{i-1}, b_i (by property (2) of D_Γ in [Theorem 3.3](#)). Applying [Lemma 3.5](#) (with the given h and $A = \Gamma(b_{i-1})$) gives the required g_i . This completes the proof of [Proposition 3.8](#). \square

References

- [1] Daniela Amato, D.Phil. Thesis, University of Oxford, 2006.
- [2] Daniela Amato, Descendants in infinite, primitive, highly arc-transitive digraphs, *Discrete Math.* 310 (2010) 2021–2036.
- [3] Daniela Amato, John K. Truss, Some constructions of highly arc-transitive digraphs, *Combinatorica* 31 (2011) 247–283.
- [4] Daniela Amato, David M. Evans, John K. Truss, Classification of some descendant-homogeneous digraphs, *Discrete Math.* 312 (2012) 911–919.
- [5] Peter J. Cameron, Oligomorphic Permutation Groups, London Math. Soc. Lecture Notes, vol. 152, Cambridge University Press, 1990.
- [6] Peter J. Cameron, Cheryl E. Praeger, Nicholas C. Wormald, Infinite highly arc transitive digraphs and universal covering digraphs, *Combinatorica* 13 (4) (1993) 377–396.
- [7] Josephine Emms, David M. Evans, Constructing continuum many countable, primitive, unbalanced digraphs, *Discrete Math.* 309 (2009) 4475–4480.
- [8] David M. Evans, Suborbits in infinite primitive permutation groups, *Bull. Lond. Math. Soc.* 33 (5) (2001) 583–590.
- [9] Clement W.H. Lam, Distance transitive digraphs, *Discrete Math.* 29 (1980) 265–274.
- [10] Rögnvaldur G. Möller, Descendants in highly arc-transitive digraphs, *Discrete Math.* 247 (1–3) (2002) 147–157.
- [11] Peter M. Neumann, Postscript to review of [5], *Bull. Lond. Math. Soc.* 24 (1992) 404–407.