# Optimal Local Inventory Management in a Decentralized Supply Chain 

Carried out at the Laboratory for Production Management and Processes, EPFL

by
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## Introduction

## Motivation of the project

This project is motivated by the course Management of Enterprise Networks and Supply Chain Management taught at École Polytechnique Fédérale de Lausanne. In this course the students have to find the best/optimal strategy for their supply chain using the XBeerGame platform. This platform simulates a 4 echelon supply chain.

The best strategy is the one that minimizes their total cost and the local cost of each echelon. That is not usually the case for a real enterprise, sometimes; the echelons are not in the same enterprise. Usually when the echelons are not part of the same enterprise usually incurs in more cost because each echelon tends to optimize their cost at expenses of the other ones.

The aim of this project is to find the optimal replenishment policy that will minimize the total cost of the supply chain.

## Research goals and approach

The goals are to find the best solution for two particular cases, so a discrete-event simulator will be created and a local search program to find the best solution. Hopefully this will serve as a basic path for future researchers that want to study how the variations of the data (demand, costs, and lead times) affect the total cost of the chain.

The approach will be creating a discrete-event simulator for each echelon that will represent the interactions over the time-steps of them. After validating the simulator, a local search program will be created in order to find the best solution.

## Thesis outline

The thesis outline is the following; first a brief literature review will be made. Then, the problem and its characteristics will be shown. After that, the models and the methodology will be explained with its assumptions and the cost model. Then the analytical cost model will be computed for a single echelon. After that, the simulation model will be code, and then the optimisation program that will find the best solution. Those programs will be programmed using C++ language.

When the results are found a sensitivity analysis will be made to validate the parameters of the optimisation program.

After the solution is found, several experiments will be done to calculate the variation of the cost caused by the variation of the service level requirement. A comparison with other local search methods will be made. The inventory and the cost of the best solution will be analysed.

After that, the two cases analysed will be compared, and discussed.

Finally some improvements and extensions will be recommended.

## Literature Review

## Meta-heuristics algorithms

An overview in meta-heuristics is studied in (Blum y Roli 2003). A framework is introduced (I\&D) in order to put different intensification and diversification components into relation with each other. This introduces an important thing to consider in the project, because there has to be a good relation of intensification and diversifications in the algorithm.

In (Keskin, Melouk y Meyer 2010) it is studied a vendor selection problem that integrates vendor selection and inventory replenishment of a firm. In this paper it is used the simulationoptimisation approach to solve the problem. It is build a discrete-event simulation model to evaluate the objective function. This approach will be used in this thesis. This works together with a scatter search-based meta-heuristic optimisation approach.

Several algorithms are used to find the best solution in an inventory management problem. In (Alrefaei y Diabat 2009) it is used a simulated annealing algorithm to solve a multi-objective inventory problem. In (Al-Rifai y Rossetti s.f.) it is solved the problem by decomposing the system by echelon and location, deriving expressions for the inventory policies parameters and developing an iterative heuristic optimization algorithm.

In (Sadeghi, y otros 2013) it is tackled the vendor-managed inventory using particle swarm optimization (PSO), in this proposed algorithm, a genetic algorithm (GA) with an improved operator is employed as a local searcher to turn it to a hybrid PSO. Also in (Yang, Chan y Kumar 2012) a GA is used to solve a single-warehouse multi-retailer replenishment system. In (Köchel y Thiem 2011) it is tackled the same problem (single-warehouse, multi-retailer system) by using PSO, that require using a simulation optimisation approach.

However, the previous papers do not tackle a multi-echelon supply chain system. Therefore, it was reviewed (Köchel y Nieländer 2005) where a simulation-optimisation algorithm was used to tackle a multi-echelon inventory system. Regarding the optimisation tool, it was used a genetic algorithm. Even if the problem is not exactly the same as in the project, that sets a starting point.

In (Zufferey 2012) a dynamic tabu search is studied applied to production. (Respen, Zufferey y Amaldi 2014) There are used several heuristics algorithms, greedy procedures, and local search methods, that will be useful to validate the tabu search on this project. This paper has a great
contribution to the thesis on the validation chapter, since it sets the basis of the validation method procedure.

## Inventory management

The literature on computing/calculating the optimal inventory policy is never ending. Publications related to the specific knowledge needed will be taken into account.

Regarding inventory management it is required to get information about inventory policies and multi-echelon management. Information about decentralized and centralized supply chain is needed as well.

In (Clark y Scarf 1960) it is tackled the multi-echelon problem analytically. This is one of the first papers that were written to face a multi-echelon supply chain. This paper uses some assumptions that will be taken in the project. For example, that the demand originates in the system at the lowest installation (retailer in the project's case) and that each echelon backlogs excess demand. It is found that the $(s, S)$ policy is optimal for the linear-plus-fixed-ordering cost case (or a fixed ordering/setup cost).

In (Lee y Whang 1999) a decentralized multi-echelon supply chain is tackled using incentives and information. This paper discusses different performance measurement involving transfer pricing, consignment, shortage reimbursement... An infinite horizon is considered. There are two echelons in the problem and the first is facing stochastic demand. Since in the project no such measures will be taken, this paper only provides a good point of view regarding the equations of backlog cost.

The aim of that problem is delivering the desired end customer service level at minimum network inventory, with the inventory divided among the various echelons. The unmet demand at the first echelon is back ordered.

A multi-echelon supply chain with decentralized control is tackled in (Axsäter, A framework for decentralized multi-echelon inventory contorl 2000). Where there is a central warehouse and a certain number of retailers. The cost structure consists in local costs and penalty cost for a delay at the warehouse. This paper has certain relevance on the project since there are used penalty costs to try to reduce the delays at the warehouse.

In (Li 2010) it is studied the inventory behaviour of autonomous and self-serving firms in a decentralized retailer-manufacturer serial supply chain. The optimal inventory is characterized analytically with and without information sharing. Some extensions are discussed in this paper:

N-stage serial systems, batch ordering policies, fixed setup costs and Markovian customer demand.

Optimal replenishment policies for all members of the supply chain are determined by hybrid meta-heuristic algorithm recently developed (Duan, Liao y Yi 2013). The hybrid meta-heuristic generates a trial solution and supplies this candidate solution to the supply chain simulation model to evaluate its objective function value. The meta-heuristic optimizer then generates a new input set according to its intelligent searching mechanism. This paper has great importance regarding this project. The same procedure will be used in the thesis.

The ( $s, S$ ) inventory policy is studied in (Baron, Berman y Perry 2010), they use the ( $s, S$ ) policy for perishable items ordered in batches. For the ( $s, Q$ ) policy there are two papers reviewed: (Rosling 2002) and (Axsäter, A simple procedure for determining order quantities under a fill rate constraint and normally distributed lead-time demand 2006).

## 1.Description of the problem

The problem consists in a decentralized supply chain. A decentralized supply chain is a chain where all the echelons work autonomously and the decision makers are distributed in each echelon. The opposite case would be a centralized supply chain, where all the decisions are taken in the headquarters.

The supply chain consists in four echelons (in downstream sense): factory, wholesaler, distributor and retailer. The product flows downstream and the information upstream. The only product commercialized is beer. The retailer is the echelon that faces the stochastic demand from the market.


Figure 1.1: Illustration of the beer distribution in the XBeerGame ${ }^{1}$
The objective is to minimize the cost of the total chain (global optimisation). The problem is complex and there are interactions between the echelons each unit of time. Therefore, a simulation-driven optimisation is chosen. Also, at it can be seen in the literature review, this is the most usual approach for problems with more than one echelons in the chain.

In principle, the XBeerGame platform is supposed to be used, but as its performance does not match the requirements, it will not be used. It is proposed to code a new simulation program in order to use it to find the total cost of the chain. That means that the simulation program will compute the objective function value that will be used in the optimisation program.

[^0]
## 2.Models and Methodology

## Chain

As said in the previous chapter, there are four echelons and two main flows in the system. 2. 1There are the information flows and the product flows.

In the Figure 2.1 is shown the Information and product flow of the chain. The red arrows represent the orders received and sent of each echelon in each time step $\left(O_{t}^{i}\right)$. That would be the information flow.

The green arrows represent the product shipment from each echelon to the downstream echelon in each time step $\left(S_{t}^{i}\right)$.

At each echelon it is represented as well the stock information that each echelon would know in a real case. The $L_{i}$ with $i=M, F, D, W, R$ is the lead time of each echelon (Manufacturer, Factory, Distributor, Wholesaler and Retailer). In this problem the manufacturer is integrated in the factory.

The on-hand inventory ${ }^{2}$ is represented by $a x_{t}^{i}$ and $x_{t}^{i}$, where the first is the initial on-hand stock at time step $t$ and the latter is the on-hand stock at the end of the time step t .

The backlogged orders ${ }^{3}$ at echelon $i$ at time step $t$ are represented by $b x_{t}^{i}$.


Figure 2.1: Scheme of the supply chain

[^1]
## Inventory policies

It has been explained the problem and the chain. But now is the time to explain the 2.2 methodology that the chain will use. As each echelon has to have a policy to manage the orders and the shipments an inventory policy has to be used.

## Event succession

An event succession has to be defined; the same will be for the two inventory policies. And the same $2 u^{1}$ ccession will be used in all the echelons.

| Time step | $t-1$ | $t+1$ | $t+2$ |
| :---: | :---: | :---: | :---: |



Figure 2.2: Scheme of the event succession
The succession of the events is the following:

1. Echelon $i$ receives the shipment from the previous echelon
a. On hand inventory $\left(a x_{t}\right)=x_{t-1}+S_{t-L}^{i-1}$
2. Echelon $i$ sends a shipment to the following echelon
a. Final inventory $\left(x_{t}\right)=a x_{t}-S_{t}^{i}$

These steps happen when there is enough quantity to deliver on time. In that case, the shipment received in every echelon is the shipment due at time $t$, it is shown in the sub index of $S_{t-L}^{i-1}$. It can be that the shipment is bigger than the order received $\left(S_{t-L}^{i-1}>O_{t}^{i-1}\right)$, this means that at this period backlogged quantity has been supplied.

In the case that there is not enough quantity to deliver the order, that quantity is backlogged $\left(b x=b x+O_{t}\right)$. The backlogged quantity is supplied as soon as there are enough pieces.

To manage the orders placement, two different inventory policies will be studied.

## 2. Models and Methodology

## Reorder point method

The reorder point method it also named $(s, Q)$ in this project. According to this policy, there are two parameters to make a decision regarding the orders.
2.2.2

In this method the orders are placed when the inventory position ${ }^{4}$ drops below $s$. So $s$ is the reorder point. As in this method the order quantity is fixed, the echelons will order a quantity $Q$. On the graph below (Figure 2.3), the lead time is 3 days, the daily demand is a normal distribution with a mean of 100 pieces and a std. dev. of 5, the reorder quantity is 400 .


Figure 2.3: Reorder point method ( $s, Q$ )
2.2.3
( $\mathrm{s}, \mathrm{S}$ ) policy
This is the other policy that will be studied in this project, also named for now on as $(s, S)$ policy. As the previous policy there are two decision parameters. Orders are placed whenever the inventory position drops below $s$ (as the $(s, Q)$ policy). But the order quantity is variable, as it is ordered the difference between the inventory position and the upper level $(S)$. The demand is the same as in the ( $s, Q$ ) case, the low level $s$ is 300 , and the upper level $S$ is 600 pieces. It can be observed that that leads to more variation on the on-hand stock.

[^2]

Figure 2.4) at time 1 and 2 it is ordered the difference between $S$ and Inventory position.


$$
\angle . L .4
$$

Figure 2.4: ( $\mathrm{s}, \mathrm{S}$ ) policy

## Assumptions

In order to be able to simulate the processes some assumptions were made. The assumptions were either to make the simulation similar to the one in the XBeerGame or to compute the cost in a straight-forward way.

Regarding the inventory management:

## 2. Models and Methodology

- There are not partial deliveries, nor missed orders.
- The same type of replenishment system is used, but, obviously, with different parameters.
- The lead time of the pieces manufacturing is constant, and the factory has unlimited capacity.
- That means that it delivers on-time always all the orders placed by the factory.
- The lead time (transportation time) is constant for all echelons.
- That does not mean that all the pieces are supplied when they are ordered.


## Cost model

2.3

The aim of the project is to find a solution with the minimum total cost. In order to do that, the cost model has to be defined.

As the total cost will be computed, the money trespassed to the echelons from other echelons has a balance of zero. These are the costs of buying pieces, the cost that comes from the price of the parts.

The fixed cost will not be taken into account, these are the costs that are independent of the volume sold or bought.

The manufacturing cost is not taken into account as it would be similar in all the cases as it relates directly to the pieces that the market demands.

The costs taken into account are:

- Launching cost $\left(\right.$ LCost $\left._{i}\right)$ : cost of placing an order
- Backlog cost ( $\mathrm{BCost}_{\mathrm{i}}$ ): cost of having an order backlogged
- Holding cost ( $\mathrm{HCost}_{\mathrm{i}}$ ): cost of having pieces in the storage (inventory)

The holding cost will be computed using the stock on-hand at the end of the period, for the back orders as well. The backorder cost will be computed with the backorders accumulated at the end of the period.

In the following table (Table 2.1) the unit costs are shown. The costs were used in the calibration of the XBeerGame during the course.

## Factory Distributor Wholesaler Retailer

| Ordering cost | $[\mathrm{m} . \mathrm{u}]$. | 25 | 25 | 25 | 25 |
| :--- | :--- | :---: | :---: | :---: | :---: |
| Inventory cost | $\left[\frac{\text { m.u. }}{\text { pieces.day }}\right]$ | 0.2 | 0.4 | 0.6 | 0.8 |
| Back order cost | $\left[\frac{\text { m.u. }}{\text { pieces.day }}\right]$ | 0.2 | 0.3 | 0.6 | 1 |

So the total cost is computed in the following form (Eq. 2.1).

$$
\begin{equation*}
\text { Total cost }=\sum_{i=1}^{4} L \operatorname{Cost}_{i}+\sum_{i=1}^{4} B \operatorname{Cost}_{i}+\sum_{i=1}^{4} H \operatorname{Cost}_{i} \tag{Eq.2.1}
\end{equation*}
$$

The demand is constant and it follows a normal distribution with a mean of 3000pcs. and a standard deviation ( $\sigma_{s}$ ) of 150pcs. To the retailer arrive 3 different orders during one day.

## 3.Analytical model of the cost

According to the literature the optimum value for a single echelon for the $(s, Q)$ policy can be found. That is assuming there are not backlog costs. This is interesting in the sense that a good initial solution can be computed to initialize the optimisation problem. This is the so called Economic Order Quantity method. That computes the optimum quantity Q balancing que holding cost and launching costs.

The total cost (Eq. 3.1) for one echelon for the entire horizon is:

$$
\begin{equation*}
C_{\text {total }}=N p+\frac{Q}{2} h+\frac{N}{Q} L \tag{Eq.3.1}
\end{equation*}
$$

Where:
N : Total demand per period $\left[\frac{\text { pieces }}{\text { t.s. } \mathbf{u}^{5} .}\right]$
p : Unit price [m.u. ${ }^{6}$ ]
Q: Order quantity [pieces]
h: Holding cost $\left[\frac{\text { m.u. }}{\text { t.s.u.pieces }}\right]$
L: Launching cost [m.u.]

Deriving the (Eq. 3.1) to find the minimum, the optimum $Q$ is found (Eq. 3.2):

$$
\begin{equation*}
Q=\sqrt{\frac{2 N L}{h}} \tag{Eq.3.2}
\end{equation*}
$$

According to the Reorder Point Method, the reorder point $s$ can be found using the equation (Eq. 3.3):

$$
\begin{equation*}
s=\text { SafetyStock }+L \cdot D \tag{Eq.3.3}
\end{equation*}
$$

Where:
L: Lead time [t. s. u.]
D: Demand in one time step $\left[\frac{\text { pieces }}{\text { t.s.u. }}\right]$
In the following table (Table 3.1) the optimum quantity is computed using (Eq. 3.2) and (Eq. 3.3). It is computed assuming their demand is constant and also is constant the lead time of the orders. The reorder point is computed as well but without taking into account the safety stock.

[^3]
## 3. AnAlytical model of the cost

The lead time of every echelon is the lead time that they have to receive the orders. That takes into account the transportation time and the processing delay of the order (1 day). For example the lead time for the wholesaler is 4 days ( 3 days transportation from the distributor +1 day of delay).

Table 3.1: Optimum quantities and reorder points for the echelons

|  | Table 3.1: Optimum quantities and reorder points for the echelons |  |  |  |  |
| :--- | :---: | ---: | ---: | ---: | ---: | ---: |
|  | Factory | Distributor | Wholesaler | Retailer |  |
| Lead time | [days] | 4 | 5 | 4 | 3 |
| Q | [pieces] | 866 | 612 | 500 | 433 |
| s | [pieces] | 12000 | 15000 | 12000 | 9000 |

## 4.Simulation model ( $\mathrm{s}, \mathrm{Q}$ ) and ( $\mathrm{s}, \mathrm{S}$ )

## Description

A simulation model is coded to compute the cost of all the echelons and to simulate their 4. Tinteractions. The simulation will follow the steps of the events as explained in the chapter 2.2.1 Event succession. In this chapter the material flow is shown. The information flow (I) and material flow (M) together follow this process:

1. Shipment received (M)
2. On-hand inventory updated (I)\&(M)
3. Order received (I)
4. Shipment sent (M)
5. On-hand inventory at the end of the period is computed $(I) \&(M)$
6. Orders are placed (I)

It is important to notice that there is a main function that simulates one period of time with this process. This function is repeated 4 times to simulate the 4 echelons and then $t_{\text {end }}$ times to simulate for echelons over a period of $t_{E n d}$.

The shipments are sent with backorders as a priority, and if they are fully satisfied, the newly received orders. So the orders are accumulated and are satisfied following a $\mathrm{FIFO}^{7}$ queue.

As it is needed to compute the service level of the retailer, the quantity sent on-time is computed at each time step.

The difference between the code of $(s, S)$ and $(s, Q)$ is only when the orders have to be placed. The initial inventory for both methods is the same and it is shown on Table 4.1. This initial inventory used is one of a calibration of the XBeerGame used during the course. As it goes downstream, the value is higher due to the stochasticity that may face the retailer. It is important to say as well, that this value can be changed, as it does not affect in a long term cost, as the on-hand value gets stable around the same value no matter the value of the initial inventory.

[^4]
## 4. Simulation model ( $\mathrm{s}, \mathrm{Q}$ ) and ( $\mathrm{s}, \mathrm{S}$ )

Table 4.1: Initial stock at each echelon

| Table 4.1: Initial stock at each echelon |  |  |  |  |  |
| :--- | :--- | ---: | ---: | ---: | ---: |
|  |  | Factory | Distributor | Wholesaler | Retailer |
| Initial inventory | [pieces] | 10000 | 12000 | 14000 | 16000 |

To simulate the demand, it is divided in 3 orders a day, the total demand is $D_{T}$ (Eq. 4.1) and the three orders are $D_{1}, D_{2}$ and $D_{3}$. The demand $D_{1}$ is the minimum of $D_{T}$ and a normal distribution of $\mathrm{N}(1000,100)$ (Eq. 4.2), that is just in case that $\mathrm{N}(1000,100)$ is greater than the $D_{T}$. So the demand $D_{1}$ will be $D_{T}$ and the others 0 (Eq. 4.3) and (Eq. 4.4). That is to make sure that $D_{1}+D_{2}+D_{3}=D_{T}$.
$D_{T}=\mathrm{N}(3000,150)$
$D_{1}=\operatorname{MIN}\left\{\mathrm{D}_{\mathrm{T}}, \mathrm{N}(1000,100)\right\}$
$D_{2}=\operatorname{MIN}\left\{D_{T}-D_{1}, \mathrm{~N}(1000,100)\right\}$
$D_{3}=\operatorname{MAX}\left\{D_{T}-D_{2}-D_{1}, 0\right\}$

## Pseudocode

4.2

The notation followed to write the pseudocode is at Table 4.2.

| $t$ | Time step $\boldsymbol{t}$ |
| :---: | :---: |
| endt | Duration of the simulation |
| $\mathrm{O}_{\mathrm{t}}^{\mathrm{i}}$ : | Orders placed by echelon i-1 to echelon $\boldsymbol{i}$ at time $\boldsymbol{t}$ |
| $\mathrm{S}_{\mathrm{t}}^{\mathrm{i}}$ : | Shipment sent from echelon $\boldsymbol{i}$ to echelon $\boldsymbol{i} \mathbf{- 1}$ at time $\boldsymbol{t}$ |
| $\mathrm{L}_{\mathrm{i}}$ : | Transport lead time from echelon $\boldsymbol{i}$ to echelon $\boldsymbol{i}$ - $\mathbf{1}$ |
| $\mathrm{x}_{\mathrm{t}}^{\mathrm{i}}$ : | Inventory level at echelon $\boldsymbol{i}$ at the end of the period $\boldsymbol{t}$ |
| ax ${ }_{\text {i }}$ : | Available inv. at echelon $\boldsymbol{i}$ at the beginning of the period $\boldsymbol{t}$ |
| bx ${ }^{\text {i }}$ | Backlog vector of echelon $\boldsymbol{i}$ following a FIFO queue |
| $Q_{\text {ontime }}$ | Cumulated quantity delivered on time from the retailer (echelon $\boldsymbol{i}=\mathbf{1}$ ) to the market ( $\boldsymbol{i}=\mathbf{0}$ ) |
| $\mathrm{Sup}_{i}$ | Replenishment upper level of ( $\boldsymbol{s}, \boldsymbol{S}$ ) policy of echelon $\boldsymbol{i}$ |
| $\mathrm{s}_{\mathrm{i}}$ | Reorder point of ( $\boldsymbol{s}, \boldsymbol{S}$ ) and ( $\boldsymbol{s}, \boldsymbol{Q}$ ) policy of echelon $\boldsymbol{i}$ |
| $\mathrm{Q}_{\mathrm{i}}$ | Order quantity of ( $\boldsymbol{s}, \boldsymbol{Q}$ ) policy of echelon $\boldsymbol{i}$ |
| $\mathrm{h}_{\mathrm{i}}$ | Inventory holding cost at echelon $\boldsymbol{i}\left[\frac{m . u}{\text { piece.t.s.u. }}\right]$ |
| $\mathrm{HCost}_{\mathrm{i}}$ | Cumulated inventory holding cost at echelon i [m.u.] |
| $\mathrm{b}_{\mathrm{i}}$ | Backorders cost at echelon i $\left[\frac{\text { m.u. }}{\text { piece-t.s.u. }}\right]$ |
| BCost $_{\text {i }}$ | Cumulated backorder cost at echelon $\boldsymbol{i}[m . u$.] |

## 4. Simulation model ( $s, Q$ ) and ( $s, S$ )

$\mathrm{l}_{\mathrm{i}} \quad$ Launching cost at echelon $\boldsymbol{i}\left[\frac{\text { m.u. }}{p i e c e}\right]$
LCost $_{\mathrm{i}} \quad$ Cumulated launching cost at echelon $\boldsymbol{i}$ [m.u.]
Suma(x) This function sums up all the elements of the vector $x$

## Initialization ( $\mathbf{t}=\mathbf{0}$ )

$i=r, w, d, f$
The initial orders are 0 for all echelons
The initial shipments are 0 for all echelons
The initial backorders are 0 for all echelons
The initial inventory at the echelons is set as the values in Table 4.1.
While ( $\mathrm{t} \leq$ endt), do:

## For each echelon do:

The on-hand inventory is updated with the shipments due today

- Set $\mathrm{ax}_{\mathrm{i}}[\mathrm{t}]=\mathrm{x}_{\mathrm{i}}[\mathrm{t}-1]+\mathrm{S}_{\mathrm{i}-1}\left[\mathrm{t}-\mathrm{L}_{\mathrm{i}-1}\right]$
> If the orders and all the backorders can be fulfilled do:
o Set $\mathrm{S}_{\mathrm{i}}[\mathrm{t}]=\mathrm{o}_{\mathrm{i}}[\mathrm{t}]+\operatorname{suma}\left(\mathrm{bx}_{\mathrm{i}}\right)$
- $\operatorname{Set} \mathrm{x}_{\mathrm{i}}[\mathrm{t}]=\mathrm{ax}_{\mathrm{i}}[\mathrm{t}]-\left(\mathrm{O}_{\mathrm{i}}[\mathrm{t}]+\operatorname{suma}\left(\mathrm{bx}_{\mathrm{i}}\right)\right)$
- Clear vector $\mathrm{bx}_{\mathrm{i}}$
- Set $Q_{\text {ontime }}=Q_{\text {ontime }}+O_{i}[t]$
> Else, if they cannot be fully fulfilled
o Satisfy the backorders in a FIFO order
- Compute $\mathrm{S}_{\mathrm{i}}[\mathrm{t}], \mathrm{x}_{\mathrm{i}}[\mathrm{t}], \mathrm{bx}_{\mathrm{i}}$ and $\mathrm{Q}_{\text {on time }}$

The new orders are placed

- Set $\operatorname{InvPos}_{i}:=$ on - hand inventory + standing orders - back orders
$\left.>\underline{\text { If }(\operatorname{InvPos}}{ }_{i}<\mathrm{s}_{\mathrm{i}}\right)$
- Set $\mathrm{O}_{\mathrm{i}-1}[\mathrm{t}]=\mathrm{Q}_{\mathrm{i}}$ (case $(\mathbf{s}, \mathbf{Q})$ )
- Set $\mathrm{O}_{\mathrm{i}-1}[\mathrm{t}]=\operatorname{Sup}_{\mathrm{i}}-\operatorname{InvPos}_{\mathrm{i}}(\boldsymbol{\operatorname { c a s e }}(\mathbf{s}, \mathbf{S}))$
- Set $_{\text {Lcost }}^{i}=$ Lcost $_{i}+l_{i}$
$>$ Else
- Set $\mathrm{O}_{\mathrm{i}-1}[\mathrm{t}]=0$

The cost at the end of the period is computed

- Set HCost $_{\mathrm{i}}=$ HCost $_{\mathrm{i}}+\mathrm{h}_{\mathrm{i}}\left(\mathrm{x}_{\mathrm{i}}[\mathrm{t}]\right)$
- Set $B \operatorname{Cost}_{i}=B \operatorname{Cost}_{i}+b_{i} \cdot \operatorname{suma}\left(b x_{i}\right)$


## 4. Simulation model ( $\mathrm{s}, \mathrm{Q}$ ) and ( $\mathrm{s}, \mathrm{S}$ )

- Set $\mathrm{S}_{0}[\mathrm{t}]=\mathrm{O}_{0}[\mathrm{t}]^{8}$
- $\mathrm{t}=\mathrm{t}+1$


## Final updates:

- Set ServiceLevel $=\frac{\mathrm{Q}_{\text {ont time }}}{\text { Accumulated demand at retailer/s }}$
- Set Total cost $=\sum$ BCost $_{i}+\sum$ LCost $_{i}+\sum$ HCost $_{i}$


## Sensitivity analysis: simulation duration

4. A. ${ }^{\text {sensitivity analysis over the simulation duration is performed. As is important to have }}$ reliable results with one run of the simulation, 10 runs are performed for different cases. So the simulation was launched 10 times over a certain periods of time ( $t=100,200,300,500,750,1000$ days). In order to decide when to stop the simulation, two graphs are made: with the service level over time and the cost over time. The cost is divided over the period of time, so they can be compared among them.

For one case ( $[s, Q]=\{10592,11992,15200,11976,361,539,615,916\}$ ) 6 different end times and 10 runs each one were made. The values of the case represent the 4 reorder points of the 4 echelons (retailer to factory order) and the 4 order quantities. It can be notice that as longer the simulation is the less variance there is (Figure 4.1 and Figure 4.2).


Figure 4.1: Graph of the daily cost over different simulation lengths

[^5]
## 4. Simulation model ( $s, Q$ ) and ( $s, S$ )



Figure 4.2: Graph of the service level over different simulation lengths
A table with the main statistics is made for this case to these observations can be corroborated with data. The confidence interval for the mean (C.I.of 95\%) is computed with the distribution t-student, as the real standard deviation behind the data is unknown. To see how big was the confidence interval (one half); it is divided over the mean. So it can be seen that as longer the period, the smaller the confidence interval gets with respect of the mean. With exception of $t=1000$ days that is a bit greater than the previous limit times. It can happen because of the case and because it is stochastic. More cases were tested (Appendix A: Output data) to make sure that more or less with a $t=500$ days the results are reliable.

| Table 4.3: Statistical information of the runs |  |  |  |  |  |  |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: |
| $\boldsymbol{t}=$ | $\mathbf{1 0 0}$ | $\mathbf{2 0 0}$ | $\mathbf{3 0 0}$ | $\mathbf{5 0 0}$ | $\mathbf{7 5 0}$ | $\mathbf{1 0 0 0}$ |
| Mean | 3278,067 | 2875,765 | 2733,408 | 2638,61 | 2567,459 | 2539,645 |
| Std. Dev. | 110,1069 | 47,06992 | 58,88274 | 35,70989 | 30,98973 | 38,90396 |
| C.I. | 78,76574 | 33,67179 | 42,12218 | 25,54532 | 22,16871 | 27,83022 |
| C.I./mean | $2,40 \%$ | $1,17 \%$ | $1,54 \%$ | $0,97 \%$ | $0,86 \%$ | $1,10 \%$ |

## 4.4

## Computation time

It is important to know the computation time for different simulation lengths. As this is an important factor for the Tabu Search. In the following graph (Figure 4.3) certain set of $s_{i}$ and $S_{i}$ are used, the set of solutions used the does not have a great impact on computation time. As it can be seen, the computation time has a linear correlation with the simulation time. For example, the computation time of 1000 days divided in 10 parts is 0.54 s, and for 500 days is 0.27s.
4. Simulation model ( $s, Q$ ) and ( $s, S$ )


Figure 4.3: Graph of the computation time depending on the simulation length

## 5.Optimisation model ( $\mathrm{s}, \mathrm{Q}$ ) and ( $\mathrm{s}, \mathrm{S}$ )

## Description

The optimisation problem, to be resolved, various steps will be followed (Figure 5.1). First, it is 5. Ineeded to define the Objective Function and its restrictions. Then, input and the initial solution have to be calculated to be able to run the optimisation. The optimisation will be a simulation driven optimisation, that means that the objective function will be computed using the simulation. Finally, the output will be the decision variables.


Figure 5.1: Optimisation problem steps

### 5.2 Optimisation problem: objective function and restrictions

The objective is to minimize the total cost of the chain. Every cost of each echelon is weighted the same, so no priorities are set. The restrictions include the service level required and the non-negativity of the variables.

As shown in chapter 2.3 Cost model, the total cost is the following:

$$
\begin{equation*}
\text { Total cost }=\sum_{i=1}^{4} L \operatorname{Cos}_{i}+\sum_{i=1}^{4} B \operatorname{Cos}_{i}+\sum_{i=1}^{4} H \operatorname{Cos} t_{i} \tag{Eq.5.1}
\end{equation*}
$$

Then the objective function for the $(s, Q)$ model and the restrictions are:

$$
\begin{equation*}
[M I N] f\left(s_{1}, s_{2}, s_{3}, s_{4}, Q_{1}, Q_{2}, Q_{3}, Q_{4}\right)=\sum_{i=1}^{4} L \operatorname{Cos}_{i}+\sum_{i=1}^{4} B \operatorname{Cos}_{i}+\sum_{i=1}^{4} H \operatorname{Cos}_{i} \tag{Eq.5.2}
\end{equation*}
$$

Service Level $\geq 95 \%$
$s_{i}>0 \forall i$
$Q_{i}>0 \forall i$

The objective function for the $(s, S)$ model and the restrictions are the same as the ( $s, Q$ ) model with exception of (Eq. 5.8). This restriction is added to avoid ordering negative quantities since the order quantity is $S_{i}-s_{i}$.

## 5. OPTIMISATION MODEL ( $s, Q$ ) AND $(s, S)$

$$
\begin{equation*}
[M I N] f\left(s_{1}, s_{2}, s_{3}, s_{4}, S_{1}, S_{2}, S_{3}, S_{4}\right)=\sum_{i=1}^{4} L \operatorname{Cost}_{i}+\sum_{i=1}^{4} B \operatorname{Cost}_{i}+\sum_{i=1}^{4} H \operatorname{Cost}_{i} \tag{Eq.5.6}
\end{equation*}
$$

Service Level $\geq 95 \%$
$s_{i}<S_{i} \forall i$
$s_{i}>0 \forall i$
$S_{i}>0 \forall i$

As said before the decision variables are $s_{i}$ and $Q_{i}$ for the $(\mathrm{s}, \mathrm{Q})$ model. For the $(\mathrm{s}, \mathrm{S})$ model $s_{i}$ and $S_{i}$ are the decision variables. With this variables the system can compute when to order and when not to.

## Optimisation program

5.3The general procedure of the tabu search will be explained in this section, and the pseudocode will take part in another chapter. The main idea of the tabu search is that it is a local search, but with forbidden moves. That means that during a certain amount of iterations it cannot go back to already evaluated solutions. This way the candidate solution does not get stuck in a local minimum. Tabu search combines intensification and diversification, regarding the intensification, the tabu search looks for the best neighbour of the best neighbour, until a local minimum. The diversification is accomplished forbidding some movements to avoid getting stuck in a local minimum.

### 5.3.1 Neighbourhood

The tabu search starts searching the neighbours of the initial solution. The variables are found in a vector: $\left[s_{1}, s_{2}, s_{3}, s_{4}, Q_{1}, Q_{2}, Q_{3}, Q_{4}\right]$ or $\left[s_{1}, s_{2}, s_{3}, s_{4}, S_{1}, S_{2}, S_{3}, S_{4}\right]$ for the (s,S) case. The maximum neighbour quantity that can be found is 16 . The first 8 neighbours come from increasing each decision variable by a certain amount. The others 8 neighbours come from decreasing the initial solution by a certain amount. The same amount is used for increasing and decreasing the value of the variables. The increasing amount and the decreasing amount are: $\Delta s$ and $\Delta Q$ or $\Delta S$ in case of the $(s, S)$ case. If a move is forbidden, the neighbour that comes from ${ }^{5}{ }^{i t} ; \mathrm{it}^{2}$ is not calculated, hence it does not get included in the neighbourhood.

## Selection of the best neighbours

The value of the objective function and the service level of the retailer are computed for every neighbour in the neighbourhood. The neighbours with a service level smaller than $95 \%$ are deleted of the possible solutions. Then, the half better neighbours are selected.

In order to decide which one is the best, the solution of every neighbour is run 10 times. This makes the tabu search more robust.

## 5. OPTIMISATION MODEL ( $s, Q$ ) AND $(s, S)$

If the best of the neighbourhood is better than the best solution found, the latter is substituted by the best neighbour.

After that, on the basis of the best neighbour, the program starts again looking for the neighbours and etc.

## Dynamic search

The tabu search looks for the solution adding or dropping a certain value to the neighbours as explained before. As this might miss better neighbours that can be reached by a smaller value of these parameters, a dynamic tabu search is used. The dynamic tabu search work as follows: when after a certain number of iterations any better solution is found, the parameters ( $\Delta s$ and $\Delta Q$ or $\Delta S$ ) are reduced by a factor X . If a better solution is found, the parameters $\Delta s$ and $\Delta Q$ or $\Delta S$ get reset to the original values.

## Initial solution

5.4

An initial solution is a very important part of the optimisation problem, because as better it is, the faster the program will found good solutions. As shown before, the optimum values for a single echelon can be found, so those values can be used to start the program.

The program saves the initial solution as the better solution ever found and then stars looking for its neighbours. In order to save a reliable value of the initial solution, it is run 10 times and then the value of the Objective function is stored. The value of 10 runs is chosen because as seen in the validation of the simulation length, within 10 runs, reliable results are found, without a great variance and a small confidence interval.
5.4.1

## [s,Q] model initial solution

For this model, the initial solution is as mentioned before; in the following table (Table 5.1) the quantities Q are shown.

Table 5.1: Initial order quantity $Q$

|  |  | Factory | Distributor | Wholesaler | Retailer |
| :--- | ---: | ---: | ---: | ---: | ---: |
| Initial $\boldsymbol{Q}_{\boldsymbol{i}}$ | [pieces] | 866 | 612 | 500 | 433 |

For the reorder point the values depend on a safety stock. So as it is really difficult to set a safety stock close to the optimum, it is set bigger. The values of the $s$ on the Table 5.2 do not take into account the safety stock.

## 5. OPtIMISATION MODEL (s,Q) AND (s,S)

| $\boldsymbol{s}$ | [pieces] | 12000 | 15000 | 12000 | 9000 |
| :--- | :--- | :--- | :--- | :--- | :--- |

There is another thing important of choosing the $s$ value: if the initial solution does not have any neighbour with a service level greater or equal than $95 \%$. The program cannot start iterating because any of these are feasible solutions. This is why is important to set a greater value of the safety stock. After many trials the following values were chosen (Table 5.3). That is setting 1000 pieces as a safety stock. As the reorder point method takes into account a constant lead time. This safety stock accounts for the variations in the lead time and variations of the demand.

Table 5.3: Initial reordering points $s_{i}$

|  |  | Factory |  | Distributor | Wholesaler |
| :--- | ---: | ---: | ---: | ---: | ---: |
|  | Retailer |  |  |  |  |
| Initial $\boldsymbol{s}_{\boldsymbol{i}}$ | [pieces] | 13000 | 16000 | 13000 | 10000 |

## [s,S] model initial solution

5.4.2

This model is a bit more complicated in terms of defining a feasible initial solution. The path taken was simulating the initial solution of the model $[\mathrm{s}, \mathrm{Q}$ ] and store the value of the inventory position at each time step. So after calculating the mean and rounded it to the thousands, the result is shown in the following table (Table 5.4). As the values are different for the different levels and s, and S, when they are rounded, the same "rounded" value is found. And after that, it was checked that the first solution and neighbours had a service level $\geq 95 \%$.

| Table 5.4: Initial $s_{i}$ and $S_{i}$ values |  |  |  |  |  |  |
| :--- | :--- | ---: | ---: | ---: | ---: | :---: |
|  |  | Factory | Distributor | Wholesaler | Retailer |  |
| Initial $\boldsymbol{s}_{\boldsymbol{i}}$ | [pieces] | 13000 | 16000 | 13000 | 11000 |  |
| Initial $\boldsymbol{S}_{\boldsymbol{i}}$ | [pieces] | 13000 | 16000 | 13000 | 11000 |  |

5.5

## Pseudocode

The notation followed to write the pseudocode is at Table 5.5.
Table 5.5: Notation of the pseudocode II

| t | Time step $\boldsymbol{t}$ |
| :--- | :--- |
| Scurrent | Current solution |
| Scand | Candidate neighbour solution |
| Sneighbour | Best candidate neighbour solution |
| Neighbours | All non-forbidden neighbours of a solution |
| Sbestever | Best solution ever found $\left(\mathrm{S}^{*}\right)$ |


| Fbestever | Objective function value of the best solution ever found (F*) |
| :--- | :--- |
| dec | Reducing factor of $\Delta s, \Delta S$ and $\Delta Q$, dec $\in[0.75,0.95]$ |
| Forbidden[i] | Vector with all the 16 moves possible, in each position, there is <br> runs_no_better_found |
| Counter that saves the number of iterations when no better |  |
| solutions are found |  |
| Neighbourhood() | Function that computes Neighbour given a solution S |
| Crop() | Function that computes the objective function of a vector of <br> neighbours, and eliminates the ones that do not reach the <br> minimum service level. And after that, eliminates the half worst |
| neighbours. |  |

### 5.5.1 <br> Main function

Initialization
Set Scurrent = Intitial solution
Set Sbestever = Scurrent
Set Fbestever = Objective function value of Sbestever
Set runs_no_better_found $=0$;
found = false;
Set $\mathrm{i}=0$;
while (Time limit not reached)

- if (runs_no_better_found $=0$ )
- Set $\Delta \mathrm{Q}=\Delta \mathrm{Q}_{0}$
- Set $\Delta \mathrm{s}=\Delta \mathrm{s}_{0}$
- if (runs_no_better_found $>\mathbf{T}$ )
- Set $\Delta \mathrm{s}=\Delta \mathrm{s} * \operatorname{dec}$
- Set $\Delta Q=\Delta Q * \operatorname{dec}$
- Set runs_no_better_found $=0$
- Call TABU()
- Set ITER=ITER+1

Final updates

- if a better solution is found
- runs_no_better_found $=0$;
- if a better solution is not found,
- runs_no_better_found = runs_no_better_found + 1;


## 5. OPTIMISATION MODEL ( $s, Q$ ) AND $(s, S)$

## TABU function

Initialization

Set SHeighbour = Scurrent
Set ties $=0$
Set Bestf $=\infty$
Set SLmin $=0.95$
Call Neighbourhood()

## Call Crop()

## While ( $i<$ Neighbours. size) do

- $\quad$ Set Scand $=$ Neighbours[i]
- If $(F(S c a n d)=$ Bestf $)$
- Set Ties $=$ Ties +1
- if random( ) $\leq \frac{1}{\text { Ties }}$
- Launch run(Costcand ${ }_{\text {avg }}$, SLcand $_{\text {avg }}$ )
- If(costcand ${ }_{\text {avg }}<$ Bestf and SLcand ${ }_{\text {avg }} \geq 0.95$ )
- $\quad$ Set Sneighbour $=$ Scand
- Set Sneighbour. cost $=$ Costcand
- $\quad$ Set Bestf $=$ Costcand
- Set Movement ${ }_{\text {done }}=$ Sneighbour.mov
- If $(\mathrm{F}($ Scand $)<$ Bestf $)$
- Launch run(Costcand ${ }_{\text {avg }}$, SLcand ${ }_{\text {avg }}$ )
- If( costcand ${ }_{\text {avg }}<$ Bestf and SLcand $_{\text {avg }} \geq 0.95$ )
- Set Sneighbour = Scand
- Set Sneighbour. cost $=$ Costcand
- Set Bestf $=$ Costcand
- $\quad$ Set Ties $=1$
- Set Movement ${ }_{\text {done }}=$ Sneighbour. mov
- $\quad$ Set $i=i+1$

Final updates:
Set Scurrent = Sneighbour
If (Bestf < Fbestever)

- Sbestever $=$ Scurrent
- Fbestever $=$ Bestf
- Set found = true
- Set TAB $=$ Uniform $(5,15)$

5. Optimisation model ( $s, Q$ ) AND $(s, S)$

- Set time $=$ TAB + ITER
- Set Forbidden[Inverse(movement)] = time


## 6. Tabu Search, sensitivity analysis

## Parameters of the Tabu search

There are several parameters in the Tabu search, so its values have to be decided on the bases 6.1 of a sensitivity analysis. As there are different parameters, a great value of combinations can be made. In the Table 6.1 are summarized all the parameters that can be changed in the Tabu search program. The maximum time allowed of the Tabu Search is 1 h for the $(\mathrm{s}, \mathrm{Q})$ case and 2 h for the $(\mathrm{s}, \mathrm{S})$ case. As the purpose of this project is academic, the iteration time could be much greater. But if one possible case for this is to let the students use it, it makes more sense to have a reduced time limit. So after observing how and when the curves converge, it is set one hour for the $(s, Q)$ and two hours for the $(s, S)$. Moreover, with these values, several experiments can be pursued in a reduced time space.

|  | Table 6.1: Parameters of the Tabu Search |
| :--- | :--- |
| Parameter | Description |
| $\Delta \boldsymbol{S}$ | Variation parameter for $S_{i}$ values |
| $\Delta \boldsymbol{S}$ | Variation parameter for $s_{i}$ values |
| $\Delta \boldsymbol{Q}$ | Variation parameter for $Q_{i}$ values |
| Dec. Factor | Decreasing factor of the variation parameters |
| T | Maximum number of iterations allowed with no better solution found. |
| TAB | Forbidden number of iterations for a certain movement $i$. |
|  | TAB $=$ Uniform $(a, b)$ |

For each case we have a set of $\Delta S / \Delta Q, \Delta s$, Dec. Factor, $T, a$ and $b$. In total, there are 6 parameters, so for example, if a set of 2 values per each parameter was implemented, there would be $2^{6}=64$ combinations.

The sensitivity analysis will be performed varying $\Delta S / \Delta Q, \Delta s$ and the decreasing factor. The other values are decided taking into account the space of solutions and number of neighbours that there are.

It is decided to set $a=5$ and $b=15$. If at maximum there are 16 neighbours, and in the case that there were tried all of the possible movements; after the $16^{\text {th }}$ iteration there would not be any possible movement if they were forbidden for more than 16 iterations. So it is decided to forbid the movements randomly between 5 and 15 iterations.

## 6. Tabu Search, sensitivity analysis

The $T$ value is set in 50 iterations. That means that if a solution is not found in 50 iterations, the values are reduced by a decreasing factor. As there are 8 decision variables, if two values were tried, that would be 265 iterations $\left(2^{8}\right)$. If those two values were the minimum and the maximum allowed of each variable, which would mean that the border of the solution space would be evaluated. So the $20 \%$ of that would be around 50 iterations. Also, there were tried different $T$ values, 20 and 100; 50 seemed to be the best fit. Even though there were not much differences, since after a certain value the behaviour was similar. Of course, other values could be used and a sensitivity analysis performed, but in this project only $\Delta S / \Delta Q, \Delta s$ and the decreasing factor are analysed.

$$
(s, Q) \text { case }
$$

For the reorder point method case ( $\mathrm{s}, \mathrm{Q}$ ), the parameters $\Delta Q, \Delta s$ and the decreasing factor are analysed. The variation parameters are $1 \%, 5 \%$ and $10 \%$ of the initial solution values; rounding it, the pairs of values are shown on the Table 6.2. The values of the decreasing factor are: $1,0.95,0.90,0.85,0.80,0.75$ and 0.70 .

|  | Table 6.2: Variation parameters |  |  |  |
| ---: | ---: | ---: | ---: | :---: |
|  | $\mathbf{1 \%}$ | $\mathbf{5 \%}$ | $\mathbf{1 0 \%}$ |  |
| $\Delta \boldsymbol{s}_{\mathbf{0}}$ | 100 | 500 | 1000 |  |
| $\Delta \boldsymbol{Q}_{\mathbf{0}}$ | 5 | 25 | 50 |  |

In the following graphs (Figure 6.1, Figure 6.2 and Figure 6.3) ${ }^{9}$ there are shown the iteration processes of Tabu search for one run of different combination of parameters.


Figure 6.1: Iteration graph of the combinations of D.f. ${ }^{10}$ with $\Delta s=100, \Delta Q=5$

[^6]
## 6. TAbu Search, sensitivity analysis



Figure 6.2: Iteration graph of the combinations of D.f. with $\Delta s=500, \Delta Q=25$


Figure 6.3: Iteration graph of the combinations of D.f. with $\Delta s=1000, \Delta Q=50$
It can be observed that as the variation parameters increase, better solutions are found faster. In the first 5 min , the searches with parameters $[\Delta s=500, \Delta Q=25]$ and $[\Delta s=1000, \Delta Q=50]$ find good solutions, while the search with $[\Delta s=100, \Delta Q=5]$ is still half way there.

But, in contrast, as the variation parameters increase, the best solution found also increases. As the variation parameters increase, the time when they found the same O.F. value is reduced, but the best solution ever found increases.

## 6. Tabu Search, sensitivity analysis

In the following table, there can be seen the best solution ever found and its time.

|  |  |  | Table 6.3: Best solution found |  |  |  |  |  |  |  |
| :--- | :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $\mathbf{1}$ | $\mathbf{0 . 9 5}$ | $\mathbf{0 . 9}$ | $\mathbf{0 . 8 5}$ | $\mathbf{0 . 8}$ | $\mathbf{0 . 7 5}$ | $\mathbf{0 . 7}$ |  |  |
| $\Delta \mathrm{s}=100$ | $\Delta \mathrm{Q}=5$ | 2466.59 | 2458.83 | 2444.14 | 2475.11 | 2484.09 | 2451.64 | 2435.01 |  |  |
| $\Delta \mathrm{~s}=500$ | $\Delta \mathrm{Q}=25$ | 2497.92 | 2511.32 | 2522.08 | 2482.76 | 2476.05 | 2493.65 | 2520.13 |  |  |
| $\Delta \mathrm{~s}=1000$ | $\Delta \mathrm{Q}=50$ | 2539.19 | 2546.14 | 2540.81 | 2554.05 | 2546.71 | 2529.50 | 2540.61 |  |  |


|  |  | Table 6.4: Time [s] when the best solution was found |  |  |  |  |  |  |  |
| :--- | :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: | :---: |
|  |  | $\mathbf{1}$ | $\mathbf{0 . 9 5}$ | $\mathbf{0 . 9}$ | $\mathbf{0 . 8 5}$ | $\mathbf{0 . 8}$ | $\mathbf{0 . 7 5}$ | $\mathbf{0 . 7}$ |  |
| $\Delta \mathrm{s}=100$ | $\Delta \mathrm{Q}=5$ | 2931.11 | 3504.65 | 3280.35 | 1869.06 | 3179.45 | 3491.27 | 3462.20 |  |
| $\Delta \mathrm{~s}=500$ | $\Delta \mathrm{Q}=25$ | 946.21 | 605.70 | 344.88 | 1904.99 | 503.40 | 706.38 | 684.35 |  |
| $\Delta \mathrm{~s}=1000$ | $\Delta \mathrm{Q}=50$ | 176.28 | 245.90 | 3323.13 | 748.32 | 2868.67 | 279.97 | 265.35 |  |

### 6.1.1.1 Validation

A validation of the Tabu Search has to be performed. Also, it is needed to validate the dynamic search with respect a static search. To validate the Tabu search it is used a Descent Local Search (DLS). For this method it is used the same code as for the Tabu Search with some differences. All moves are allowed in all steps, the initial solution is the initial solution of the tab search with a perturbation applied at all components of the vector. The perturbation is a random value $z \in[0.75,1.5]$. After a certain amount of runs (50) without finding any better solution, the search restarts with a different initial solution, until the end time is reached.

The validation of the dynamic tabu search is performed with a static tabu search, with all the parameters as the dynamic tabu search but the decreasing factor. The decreasing factor is 1 as it is a static search.

Table 6.5: Objective function value of the best solution found at each run

| Dynamic TS | Static TS | DLS |
| ---: | ---: | ---: |
| $F^{*}[\mathrm{~m} . \mathrm{u}]$. | $\mathrm{F}^{*}[\mathrm{~m} . \mathrm{u}]$. | $\mathrm{F}^{*}[\mathrm{~m} . \mathrm{u}]$. |
| 2511.6 | 2465.7 | 2501.5 |
| 2500.0 | 2491.9 | 2750.1 |
| 2473.6 | 2403.6 | 2568.5 |
| 2473.3 | 2451.8 | 2529.9 |
| 2475.6 | 2432.6 | 2786.9 |
| 2445.7 | 2431.2 | 3158.6 |
| 2474.4 | 2441.1 | 2821.7 |
| 2482.4 | 2427.9 | 2497.1 |
| 2462.7 | 2436.7 | 3222.3 |
| 2473.3 | 2435.5 | 2492.3 |

## 6. TABU SEARCH, SENSItivity ANALYSIS

Table 6.6: Time passed until the best solution is found at each run

| Dynamic TS |  | Static TS |
| ---: | :---: | ---: |
| $\mathrm{t}^{*}[\mathrm{~s}]$ | $\mathrm{t}^{*}[\mathrm{~s}]$ | tLS |
| 2323.9 | 2754.9 | 1829.3 |
| 81.8 | 2226.4 | 2859.5 |
| 709.3 | 2728.8 | 3018.3 |
| 3531.3 | 3274.0 | 435.0 |
| 422.0 | 3489.5 | 1791.8 |
| 1905.6 | 2905.6 | 1253.4 |
| 1330.3 | 3401.8 | 1357.7 |
| 1253.7 | 3147.6 | 313.8 |
| 2361.1 | 1833.7 | 0.0 |
| 3531.3 | 2700.6 | 1407.5 |

On the previous tables (Table 6.5 and Table 6.6) the values of the several runs are shown, the best solution found, and the time passed until it was found.

In the following graphs, to observe better the difference among the three programs, a boxplot is made. So all the values found on the previous tables are represented using quartiles. Also, the confidence interval of the $95 \%$ for the mean is shown as well.

One of the better combinations is $\Delta s=500$ and $\Delta Q=25$ with a decreasing factor of 0.8 . So those will be the parameters for the DLS and TS static. So the three methods will be compared in the following graphs, showing the best solution found, and the time when it was found.


Figure 6.4: Boxplot of the best solution found in 10 runs for DTS, STS and DLS

## 6. Tabu Search, sensitivity analysis



Figure 6.5: Boxplot of the best solution found in 10 runs for DTS and STS


Figure 6.6: Boxplot of the time when the best solution was found in 10 runs for DTS, STS and DLS
In the graphs it can be seen that the descent local search does not perform very well, its time is fast but the solutions are far the best solution found using the tabu methods.

It can be concluded that for this case, even though the Static TS is faster, the dynamic tabu search finds better solutions than the STS. So the dynamic tabu search performs better for $\Delta s=500$ and $\Delta Q=25$ than the static tabu search.

## 6. TABU SEARCH, SENSItivity ANALYSIS

## ( $\mathrm{s}, \mathrm{S}$ ) case

For the reorder point method case $(s, S)$, the parameters $\Delta S, \Delta s$ and the decreasing factor are analysed. The variation parameters are $1 \%, 5 \%$ and $10 \%$ of the initial solution values; rounding 6.1.2 it, the pairs of values are shown on the Table 6.2. The values of the decreasing factor are: $1,0.95,0.90,0.85,0.80,0.75$ and 0.70 . After running the program several times, it is decided a duration of 2 h , observing he graphs it is deducted that 1 h is not enough.

|  | Table 6.7: Variation parameters |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  | $\mathbf{1 \%}$ | $\mathbf{5 \%}$ | $\mathbf{1 0 \%}$ |  |
| $\Delta \boldsymbol{s}_{\mathbf{0}}$ | 100 | 500 | 1000 |  |
| $\Delta \boldsymbol{S}_{\mathbf{0}}$ | 100 | 500 | 1000 |  |

In the following graphs (Figure 6.7, Figure 6.8 and Figure 6.9) are shown the iterations for combinations $\Delta \mathrm{s}, \Delta \mathrm{S}$ and several decreasing factor, as in the ( $\mathrm{s}, \mathrm{Q}$ ) method.


Figure 6.7: Iteration graph of the combinations of D.f. with $\Delta s=100, \Delta S=100$

## 6. Tabu Search, sensitivity analysis



Figure 6.8: Iteration graph of the combinations of D.f. with $\Delta s=500, \Delta S=500$


Figure 6.9: Iteration graph of the combinations of D.f. with $\Delta s=1000, \Delta S=1000$
On those graphs it can be observed that for great variations: 500 and 1000, at the beginning descent quicker than for the 100 variation. But after that huge descent almost all the curves go flat. That does not happen in the others graphs (Figure 6.8 and Figure 6.9) as after a while (30

## 6. Tabu Search, sensitivity analysis

min for the first and 50 for the second) they have an important descent towards the best solution.

In this method $(s, S)$ the decreasing factors have more impact as the variation parameters increase. In the first graph (Figure 6.7) there is no such correlation, but in the others (Figure 6.8 and Figure 6.9) the value of the best solution is almost sorted as the decreasing factor. That can be observed very well in Figure 6.9, where all except D.f $=0.85$ and 0.9 are sorted exactly the same as the best solution found.

To sum up, a table is made to show in numbers what it is on the graphs.

|  |  | 1 | 0.95 |  | 0.85 | 0.8 | 0.75 | 0.7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | 0.9 |  |  |  |  |  |
| $\Delta s=100$ | $\Delta \mathrm{S}=100$ |  | 3229.12 | 3311.98 | 3210.94 | 3385.22 | 3303.31 | 3210.37 | 3394.29 |
| $\Delta s=500$ | $\Delta S=500$ | 3655.59 | 3567.57 | 3482.33 | 3208.19 | 3361.08 | 3225.61 | 3409.01 |
| $\Delta s=1000$ | $\Delta S=1000$ | 3698.2 | 3684.42 | 3672.86 | 3579.33 | 3640.99 | 3487.16 | 3467.79 |


|  | Table 6.9: Time [s] when the best solution was found |  |  |  |  |  |  |  |
| :--- | :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $\mathbf{1}$ |  | $\mathbf{0 . 9 5}$ | $\mathbf{0 . 9}$ | $\mathbf{0 . 8 5}$ | $\mathbf{0 . 8}$ | $\mathbf{0 . 7 5}$ | $\mathbf{0 . 7}$ |
| $\Delta \mathrm{s}=100$ | $\Delta \mathrm{~S}=100$ | 4566.82 | 6662.46 | 6175.19 | 2813.79 | 7172.53 | 5762.40 | 5337.97 |
| $\Delta \mathrm{~s}=500$ | $\Delta \mathrm{~S}=500$ | 2112.47 | 7052.35 | 5455.32 | 7046.10 | 6247.64 | 6603.49 | 5927.50 |
| $\Delta \mathrm{~s}=1000$ | $\Delta \mathrm{~S}=1000$ | 1089.44 | 6806.67 | 4642.94 | 7178.22 | 4015.96 | 6887.45 | 7181.02 |

### 6.1.2.1 Validation

This method requires a validation as well. So as explained in the previous chapter, a Descent Local Search and a Static Tabu search are performed. The parameters will be $\Delta s=100$ and $\Delta S=100$, as this combined with 0.75 of decreasing factor in the Dynamic Tabu Search performs very well.

Table 6.10: Objective function value of the best solution found at each run II

| Dynamic TS | Static TS | DLS |
| ---: | ---: | ---: |
| $F^{*}[\mathrm{~m} . \mathrm{u}]$. | $F^{*}[\mathrm{~m} . \mathrm{u}]$. | $F^{*}$ [m.u.] |
| 3264.0 | 3212.2 | 3643.2 |
| 3224.5 | 3224.8 | 6820.4 |
| 3339.8 | 3272.4 | 3536.2 |
| 3302.9 | 3364.3 | 5487.1 |
| 3230.3 | 3367.2 | 4908.5 |
| 3327.6 | 3232.8 | 3539.7 |
| 3268.4 | 3278.5 | 3603.6 |
| 3199.0 | 3239.4 | 3696.9 |
| 3299.2 | 3323.1 | 4496.5 |
| 3319.6 | 3313.4 | 4165.3 |

## 6. TABU SEARCH, SENSItivity ANALYSIS

Table 6.11: Time passed until the best solution is found at each run II

| Dynamic TS | Static TS | DLS |
| :---: | :---: | ---: |
| $\mathrm{t}^{*}[\mathrm{~s}]$ | $\mathrm{t}^{*}[\mathrm{~s}]$ | $\mathrm{t}^{*}[\mathrm{~s}]$ |
| 6936.0 | 2819.2 | 1198.3 |
| 6074.0 | 2551.3 | 2094.5 |
| 5698.2 | 4004.3 | 260.9 |
| 7154.5 | 6772.0 | 3490.5 |
| 6252.1 | 3021.4 | 5334.8 |
| 6829.0 | 6775.2 | 1011.8 |
| 6977.1 | 6986.8 | 366.7 |
| 4285.2 | 5731.8 | 584.5 |
| 6190.6 | 6342.3 | 1279.0 |
| 7105.6 | 6785.8 | 2632.2 |

As on the previous chapter, the values of the runs can be found on a table (Table 6.10 and Table 6.11). The values are the objective function value of the best solution found at each run, and the time when it is found.

In order to compare in a more visual way the values, two boxplots are made, one with the best values found in the 10 runs and the second one with the time when the best solution is found.


Figure 6.10: Boxplot of the best solution found in 10 runs for DTS, STS and DLS
6. TABU Search, sensitivity analysis


Figure 6.11: Boxplot of the best solution found in 10 runs for DTS and STS


Figure 6.12: Boxplot of the time when the best solution was found in 10 runs for DTS, STS and DLS
The Descent Local search does not perform as the tabu search; the solutions found are far greater than the found by the two tabu methods.

The dynamic tabu search performs better than the static, for this pair of $\Delta s$ and $\Delta S$ the confidence interval for the mean is smaller. Also the extremes of the confidence interval of a 95\% are a bit below the values of the static tabu search.

## 6. Tabu Search, sensitivity analysis

The time when the dynamic tabu search founds the best solution is greater than with the Static Tabu search. That happens because the dynamic tabu search keeps decreasing the variation parameter with the factor 0.75 until it founds a better solution.

There is not as much difference as in the ( $s, Q$ ) case because the variation parameters $\Delta s=100$ and $\Delta S=100$ are not as great as in the $(s, Q)$. So the decreasing values do not affect as much.

## 7. Results

## 7.Results

## ( $s, Q$ ) case best solution

Using the dynamic tabu search, the best solution is found and shown in the following table:
7.1

| Table 7.1: Solution found for $(\mathrm{s}, \mathrm{Q})$ case |  |  |  |  |  |
| ---: | ---: | ---: | ---: | ---: | ---: |
|  |  | Retailer | Wholesaler | Distributor | Factory |
| $\boldsymbol{s}_{\boldsymbol{i}}$ | [pieces] | 10071 | 12275 | 15204 | 11884 |
| $\boldsymbol{Q}_{\boldsymbol{i}}$ | [pieces] | 420 | 560 | 644 | 966 |

The confidence interval of $95 \%$ for the daily cost is [2372.56, 2402.41] and the confidence interval for the service level is: $[0.937,0.957]$. In the following graphs there can be observed the inventory level and the on-hand stock at the end of each period for all the echelons of one simulation. In order to be able to appreciate the peaks, only 100 days are shown. It is important to remind that each day has 10 steps, so between 0 and 10 days, there are 100 steps.


Figure 7.1: Inventory level, on-hand stock at the end of each period, and reorder point for the Factory

## 7. ReSults



Figure 7.2: Inventory level, on-hand stock at the end of each period, and reorder point for the Distributor


Figure 7.3: Inventory level, on-hand stock at the end of each period, and reorder point for the Wholesaler


Figure 7.4: Inventory level, on-hand stock at the end of each period, and reorder point for the Retailer
It can be observed that the echelons order more frequently in a downstream sense. That is because the order quantity decreases as it goes downstream.

The average on-hand stock at the end of the period is shown in the following graphs. The total on-hand stock average is 3291 pieces, more than the mean of the daily demand.


Figure 7.5: Average on-hand stock at the end of the period of 500 days for each echelon

As can be seen on the previous figure, the on-hand stock at the retailers is the greatest one, the stock increases from the wholesaler to the distributor. Some factors can be the decreasing value of the holding cost. The factory does not have that great on-hand stock probably caused by the small holding cost and the reduced lead time.

## 7. Results

## Results for several service levels

It is interesting also to know how could vary the solution and the cost if the service level was lower and higher than $95 \%$. So two searches were launched to compute the best solution for 7.1.1 the case $S L=0.90$ and $S L=0.98$. As the case is stochastic, even though a solution was found for a $S L=0.98$, when the solution was simulated 20 times, the average service level was 0.97 .

The solution found are shown in the following table (Table 7.2), the solution for a $95 \%$ service level is the one found in the previous chapter.

Table 7.2: Best solution for service levels: $90 \%, 95 \%$ and $97 \%$

|  | $\boldsymbol{S}_{\boldsymbol{i}}$ |  |  |  | $\boldsymbol{Q}_{\boldsymbol{i}}$ |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | F | D | W | R | F | D | W | R |
| SL=90\% | 11614 | 15261 | 12327 | 10000 | 926 | 579 | 548 | 411 |
| SL=95\% | 11884 | 15204 | 12275 | 10071 | 966 | 644 | 560 | 420 |
| SL=97\% | 11976 | 15200 | 11992 | 10592 | 916 | 615 | 539 | 361 |

In the following graphs (Figure 7.6) are represented the values on the table. In the reorder point level a slightly tendency can be observed. But it cannot be assured that it is caused by the service level, as it could possible happen that another fitting solution could show a different tendency.


Figure 7.6: Reorder point levels for the service levels: $90 \%, 95 \%$ and $97 \%$


Figure 7.7: Order quantity for the service levels: 90\%, 95\% and 97\%
It can be observed in the following graph that there is an important difference regarding the average on-hand stock at the retailer's at it increases as the service level increases.


Figure 7.8: Average on-hand stock of the simulation using the solutions found for $\mathrm{SL}=90 \%, 95 \%$ and $97 \%$

Ten simulations were run to enclose the daily cost for the different service levels. The intervals of confidence for the mean of the daily cost were computed (Table 7.3). It can be noticed that they do not overlap. So it can be said that in $95 \%$ of cases the average daily cost of $\mathrm{SL}=90 \%$ will be smaller than the daily cost for a $\mathrm{SL}=95 \%$, and also, the daily cost of $\mathrm{SL}=95 \%$ will be smaller than the daily cost of the solution of $S L=97 \%$.

| Table 7.3: 95\% confidence interval for the mean of the daily cost |  |
| :--- | :---: |
|  | 95\% Confidence interval for the mean |
| Service level $=\mathbf{9 0} \%$ | $\{2329.50,2368.94\}$ |
| Service level $=\mathbf{9 5} \%$ | $\{2387.55,2433.03\}$ |
| Service level $=\mathbf{9 7 \%}$ | $\{2604.03,2655.93\}$ |

## 7. Results

In the following graph (Figure 7.9) it is shown the values in a boxplot of the 10 runs for each solution of service level. In the 2 Q box are shown the second quartile and in the 3 Q box the third quartile. The red cross is the mean, and it can be seen that the mean does not follow a straight line, so the daily cost is not linearly proportional to the service level.


Figure 7.9: Boxplot of the runs of the solutions for service levels $=90 \%, 95 \%$ and $97 \%$

## $7.2 \quad(\mathbf{s}, \mathbf{S})$ case best solution

The solutions for the ( $\mathrm{s}, \mathrm{S}$ ) case are shown in the following table (Table 7.4). It can be observed that the levels are really close, that shows that it is rather closer to a base stock (only one level) policy than to an ( $s, S$ ) policy. So this model could be simplified to one when each period is ordered the difference between $S$ and the inventory position, and not only when the inventory position is smaller than s.

Table 7.4 Best solution found for $(\mathrm{s}, \mathrm{S})$ case

|  |  | Factory | Distributor | Wholesaler | Retailer |
| :--- | :--- | ---: | ---: | ---: | ---: |
| $\boldsymbol{s}_{\boldsymbol{i}}$ | [pieces] | 12056 | 15940 | 12163 | 10419 |
| $\boldsymbol{S}_{\boldsymbol{i}}$ | [pieces] | 12124 | 16090 | 12178 | 10430 |

In the following graphs there are shown the inventory position, s level and on-hand stock at the end of each period. It can be notice that as more up-stream, the peaks of the inventory position increase. And that results in greater on-hand stock.


Figure 7.10: Inventory level, on-hand stock at the end of each period, and reorder point for the Factory


Figure 7.11: Inventory level, on-hand stock at the end of each period, and reorder point for the Distributor


Figure 7.12: Inventory level, on-hand stock at the end of each period, and reorder point for the Wholesaler


Figure 7.13: Inventory level, on-hand stock at the end of each period, and reorder point for the Retailer

A graph representing the average on-hand stock is made, in this graph (Figure 7.14) the quantity is shown, and it shows what appears in the previous graphs, the on-hand stock is accumulated at the distributor's and at the factory's. The addition of the average on-hand stock is: 5266 pieces.

## 7. Results



Figure 7.14: Average on-hand stock at the end of the period of 500 days for each echelon II

## Results for several service levels

As d6ine with the ( $s, Q$ ) case, the comparison between different service levels is done as well for the $(s, S)$ case. The service levels are $S L=0.88,0.94$ and 0.97 . A table with the solutions is made (Table 7.5).

Table 7.5: Best solution for service levels: $88 \%, 94 \%$ and $97 \%$

|  | $\boldsymbol{s}_{\boldsymbol{i}}$ |  |  |  | $\boldsymbol{S}_{\boldsymbol{i}}$ |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | F | D | W | R | F | D | W | R |
| SL=88\% | 11891 | 15768 | 12039 | 10345 | 15815 | 12142 | 10345 | 11891 |
| SL=94\% | 12056 | 15940 | 12163 | 10419 | 16090 | 12178 | 10430 | 12056 |
| SL=97\% | 12619 | 15195 | 12884 | 10498 | 15220 | 12896 | 10555 | 12619 |

The values of the table are represented on the two following figures (Figure 7.15 and Figure 7.16). It can be observed a slight tendency but it is not conclusive because not all the fitting solutions and for all service levels between 0.88 and 0.97 are computed.


Figure 7.15: Order point $s_{i}$ for the service levels: $88 \%, 94 \%$ and $97 \%$


Figure 7.16: Upper level $S_{i}$ for the service levels: $88 \%, 94 \%$ and $97 \%$
The on-hand stock (Figure 7.17) shows that the solution for the $97 \%$ service level is greater than for the other cases. That can be caused for the high service level, as the retailer has to accomplish a $98 \%$ and has to have enough stock to satisfy on time the demand.


Figure 7.17: Average on-hand stock of the simulation using the solutions found for $\mathrm{SL}=88 \%, 94 \%$ and $97 \%$
Ten simulations were run to enclose the daily cost for the different service levels as in the previous chapter with the ( $s, Q$ ) case. The confidence interval for the mean daily cost are computed and shown on the table (Table 7.6). The intervals do not overlap as in the previous case.

Table 7.6: 95\% confidence interval for the mean of the daily cost II
95\% Confidence interval for the mean

| Service level $=\mathbf{8 8} \%$ | $\{3112.59,3154.14\}$ |
| :--- | :--- |
| Service level $=\mathbf{9 4} \%$ | $\{3195.22,3243.36\}$ |
| Service level $=\mathbf{9 7} \%$ | $\{3401.98,3451.87\}$ |

## 7. Results

In the following graph (Figure 7.18) it is shown the values in a boxplot of the 10 runs for each solution of service level. In the 2 Q box are shown the second quartile and in the 3 Q box the third quartile. The red cross is the mean, and it can be seen that the mean does not follow a straight line, so as in the previous case the daily cost is not linearly proportional to the service level.


Figure 7.18: Boxplot of the runs of the solutions for service levels $=88 \%, 94 \%$ and $97 \%$

## 8.Discussion

## ( $\mathrm{s}, \mathrm{Q}$ ) case

In the ( $s, Q$ ) case the echelon that has the more cost is the retailer, and this cost is mostly 8.1caused by the holding cost of the on-hand inventory. In the following figure (Figure 8.1) are shown the portions of each cost over the total cost of the echelon. $B_{i}$ is the backorder cost at echelon $i, i=r, w, d, f . H_{i}$ is the holding cost at echelon $i . L_{i}$ is the launching cost at echelon $i$, this represents the cost that incurs when an order is placed, it does not depend on the quantity ordered, only on the frequency.

The retailer only has a $2 \%$ of cost caused by the backorder cost, this happens because the retailer has to face a $95 \%$ of service level, so not much backorders are allowed.

The wholesaler also has most of the cost caused by the holding cost, but almost a quarter is caused because of the launching cost. The backorder cost is a $12 \%$ of the total cost; this is because the wholesaler does not have any service level requirement. The distributor has also the same distribution as the wholesaler. But the factory has even less holding cost portion than the others.


## 8. DISCUSSION



Figure 8.1: Distribution of the cost of each echelon

The following figure shows the aggregated costs of the echelons, so each colour represents a different echelon, and each tone of the colour, the different costs (Figure 8.2).

The retailer causes the $47 \%$ of the total cost, the wholesaler the $23 \%$, the distributor the $19 \%$ and the factory the remaining $11 \%$.

To remind (Table 8.1), the average on-hand stock of each echelon, the holding cost, and the order quantity:

Table 8.1: Avg. on-hand stock, holding cost and order quantity

|  | Factory | Distributor | Wholesaler | Retailer |
| :---: | :---: | :---: | :---: | :---: |
| Avg. on-hand stock [pcs] | 763 | 751 | 594 | 1184 |
| $h_{i}\left[\frac{m . u .}{p c s \cdot d a y}\right]$ | 0.2 | 0.4 | 0.6 | 0.8 |
| $\mathrm{Q}_{i}[\mathrm{pcs}]$ | 966 | 644 | 560 | 420 |

As a simplification it could be said that the holding cost of each echelon is the avg. on-hand stock times the holding cost times the number of days. So here it can be seen that the retailer will have the greatest holding cost as is the one that has the greatest on-hand stock and holding cost. Even though the distributor and the factory have a very close average on-hand stock, as the holding cost is different, the portion over the total cost is different too. The distributor has more stock (751pcs.) than the wholesaler (594pcs.) but as the cost is greater at the latter, the cost is greater for the wholesaler than the distributor.

## 8. DISCUSSION

The total launching cost decreases in an upstream sense, this is caused because the order quantity increases in the upstream sense which means that the frequency decreases, and that leads to less cost. The launching cost per order is the same for all echelons ( $25 \mathrm{~m} . \mathrm{u}$./order).

With exception of the retailer, the backorder costs increases in a downstream sense, caused by the increase of the unitary cost.


Figure 8.2: Distribution of all the costs of each echelon
8.1.1

## Several service levels

It can be observed that as the service level increases, the inventory increases so the holding cost increases as well. Another main difference caused by the variation of the service level is that the backorder cost proportion decreases when the service level increases. The launching cost follows the same rule as the solution for $S L=0.95$ as it is still proportional to frequency of the shipments and inversely proportional to the ordering quantity.

## 8. DISCUSSION




Figure 8.3: Comparison between the cost distributions of the solutions with $\mathrm{SL}=0.92$ and $\mathrm{SL}=0.98$

To better understand the cost, a graph (Figure 8.4) with the absolute cost is made, it can be seen that at the wholesaler, distributor and factory the difference is slight. However, at the retailer's the difference of the total cost is more important than at the other echelons. The cost increases with the service level, caused by the increase of on-hand stock at the retailer's.


Figure 8.4: Daily cost of each echelon for each solution ( $\mathrm{SL}=0.92,0.95,0.98$ )
8.2

## (s,S) Case

The distribution of the cost for this policy is different than the ( $s, Q$ ) policy (Figure 8.5). The retailer still has a really small backorder cost with respect of the other costs. The wholesaler has more than a quarter of backorder cost, but the cost is still dominated by the holding cost. The distributor has a really huge proportion of holding cost caused by the variation of orders and quantity of the downstream echelons. The factory has a big holding cost, almost three quarters of the total.


Figure 8.5: Distribution of the cost of each echelon, best solution $(s, S)$ policy

The retailer causes the $37 \%$ of the total cost, the wholesaler a $23 \%$, the distributor another $23 \%$ and the remaining $17 \%$ is caused by the factory (Figure 8.6).

The launching cost decreases in an upstream sense, knowing that the unitary launching cost is the same for all the echelons, that means that the frequency increases in a downstream sense.

The distributor and the factory have a big holding cost caused by the variation and the peaks of the orders coming from the downstream echelons. That is because the $(\mathrm{s}, \mathrm{S})$ model is very reactive, the echelons react too much to the peaks, so this effect increases in an upstream sense, that effect can be observed on the graph of on-hand inventory. The peaks on the inventory position cause having more on-hand inventory.

The backorders costs are small at all the echelons with exception of the wholesaler that receives the non-constant orders of the retailers almost every time step. As the orders of the retailers are not constant the wholesaler has to face more stochasticity than the retailer's, and that leads to backorders.

## 8. DISCUSSION



Figure 8.6: Distribution of all the costs of each echelon

### 8.2.1 Several service levels

For the $(\mathrm{s}, \mathrm{S})$ case the distribution of the cost for both service levels are shown below(Figure 8.7). The proportion of the holding cost at the retailers increases. There is a shift between the wholesaler and the distributor, as the distributor reduces its holding cost and the distributor increases it.


Figure 8.7: Comparison between the cost distributions of the solutions with $\mathrm{SL}=0.88$ and $\mathrm{SL}=0.98$

As on the other policy the increasing service level causes increasing cost, it is shown in the following figure (Figure 8.8). In this particular run the cost of the solution of $\mathrm{SL}=0.88$ is higher than the solution for $S L=0.94$, this can happen, as seen in the previous chapter the confidence

## 8. DISCUSSION

interval overlap. Almost all the cost of the $\mathrm{SL}=0.88$ solution are higher than the solution of $S L=0.94$. But not at the retailer's, so it can be assumed that the mean would be below that point. So as in the other policy there is a clear tendency at the retailer's, where the cost increases as the service level.


Figure 8.8: Daily cost of each echelon for each solution ( $\mathrm{SL}=0.88,0.94,0.97$ )

## $8.3 \quad(\mathbf{s}, \mathbf{Q})$ and $(\mathbf{s}, \mathbf{S})$ case comparison

This chapter will provide a short summary of the cost results for the best solutions found for both policies.

One of the main differences that trigger the difference on the cost is the on-hand inventory. As it can be observed on the graph below, there is not much difference at the retailers and wholesaler, but at the distributor and factory the value for the $(\mathrm{s}, \mathrm{S})$ is the double of the ( $\mathrm{s}, \mathrm{Q}$ ) value.

Table 8.2: Avg. on-hand stock at each echelon for the solutions of ( $s, Q$ ) and ( $s, S$ ) policies

| On-hand stock[pcs.] | Factory | Distributor | Wholesaler | Retailer | Total |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $(\mathrm{s}, \mathrm{Q})$ solution | 762.7 | 751.1 | 593.5 | 1183.7 | 3291.0 |
| $(\mathrm{~s}, \mathrm{~S})$ solution | 1880.6 | 1438.6 | 679.4 | 1267.8 | 5266.4 |

On the figure below (Figure 8.9) it can be observed that the distribution of the cost is quite different. For the first case the holding cost at the retailers is almost $40 \%$ while on the other case it is $32 \%$. In the second policy the distributor and the factory have more share than in the $(s, Q)$ policy.


Figure 8.9: Cost distribution of the best solution of policies $(\mathrm{s}, \mathrm{Q})$ and $(\mathrm{s}, \mathrm{S})$
The absolute cost values are also different (Table 8.3), being the cost of the $(\mathrm{s}, \mathrm{S})$ greater than the cost of the $(s, Q)$, the confidence interval do not overlap by far.

Table 8.3: Cost and service level confidence intervals comparison
Total cost (95\% C.I.) Service level (95\% C.I.)

| $(s, Q)$ policy | $[2387.55,2433.03]$ m.u. | $[0.937,0.957]$ |
| :--- | :--- | :--- |
| $(s, S)$ policy | $[3195.22,3243.36]$ m.u. | $[0.937,0.954]$ |

The difference can be seen well in the following graph, where a boxplot is made and the confidence interval shown.


Figure 8.10: Boxplot of the daily cost of 10 runs for policies $(s, Q)$ and $(s, S)$

For each echelon the absolute costs are computed to show that difference at the distributor and factory is mainly caused by the holding cost (


Figure 8.11). As it can be observed, the Distributor and Factory almost double the holding cost when a $(\mathrm{s}, \mathrm{S})$ policy is used instead of a ( $\mathrm{s}, \mathrm{Q}$ ) policy. In this case, the wholesaler has more backorder cost with he $(s, S)$ policy than in the $(s, Q)$ policy.



Figure 8.11: Absolute cost at each echelon for $(s, Q)$ and $(s, S)$ policy
This differences do not mean that the $(\mathrm{s}, \mathrm{S})$ policy is worse than the $(\mathrm{s}, \mathrm{Q})$ policy. That means that in this particular case with its unit costs, lead times and demand it is less costly to use the $(s, Q)$ policy instead of a ( $s, S$ ) policy.

## 9.Improvements and extensions

## Improvements

The program can always be improved. In this case, some extensions were thought to be 9. 1implemented in the future. Regarding the simulation, an extension would be making the parameters s, Q or S depending on the demand or orders received, that would mean that the parameters would learn as the simulation goes on. That would be applicable to a decreasing or increasing demand, for example.

An improvement regarding the Tabu search program is when a tie happens. In the Tabu search of this project, the value has to be exactly the same to consider a tie. So an improvement can be to allow some difference between the best value found and the candidate. That would mean that: being $F$ the best value and $F_{\text {candidate }}$. A tie would occur when:
$0.95 \times F \leq F_{\text {candidate }} \leq 1.05 \times F_{\text {candidate }}$

A value $F_{\text {candidate }}<0.95 \times F$ would be considered directly a better value and not a tie.

### 9.2 Extensions

The previous chapters show that there is a high difference between the two policies studied. There is a big difference as well depending on the service level restriction, that as higher the service level, the higher the cost.

In this project the global cost is analysed, but an interesting extension would be to optimize only one echelon. Only the cost of this echelon would be computed on the objective function. It could be studied what would happen if a single echelon tried to optimize its cost. Would it be really different than the result of optimizing the global cost of the chain?

An interesting extension would be to analyse how changing the cost structure can impact on the solutions and on the cost distribution. Does higher holding cost reduce the on-hand stock and increase the launching cost?

Varying the lead time may have an impact on the solution and on the ordering variables, but does it imply that the cost will vary?

Other possible extension is analysing how the variation of the demand may affect the results. In this project a constant demand has been used, but it can be changed to a cyclical, seasonal

## 9. IMPROVEMENTS AND EXTENSIONS

or increasing/decreasing demand. Changing the demand may affect the differences of the ( $s, Q$ ) and $(s, S)$ best cost making the difference greater or smaller.

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## 11. Appendix A: Output data

Sensitivity analysis: Simulation duration

## Combination set of:

| s: | $\{11000$, | 12000 | 15200 | $12000\}$ |
| :---: | :---: | :---: | :---: | :---: |
| Q: | $\{400$, | 550 | 700 | $800\}$ |




## Statistical Information

|  | $\mathbf{1 0 0}$ | $\mathbf{2 0 0}$ | $\mathbf{3 0 0}$ | $\mathbf{5 0 0}$ | $\mathbf{7 5 0}$ | $\mathbf{1 0 0 0}$ |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| Mean | 3625,665 | 3261,021 | 3135,353 | 3031,532 | 2993,145 | 2951,08 |
| Std Dev | 77,20688 | 56,67508 | 60,42102 | 50,34162 | 26,2887 | 17,34614 |
| CI | 55,23048 | 40,54291 | 43,22259 | 36,01223 | 18,8058 | 12,40868 |
| CI/mean | $1,52 \%$ | $1,24 \%$ | $1,38 \%$ | $1,19 \%$ | $0,63 \%$ | $0,42 \%$ |

## Combination set of:

| s: | $\{11000$, | 11000 | 13200 | $11000\}$ |
| :---: | :---: | :---: | :---: | :---: |
| Q: | $\{400$, | 550 | 700 | $800\}$ |


11. Appendix A: Output data


## Statistical Information

|  | $\mathbf{1 0 0}$ | $\mathbf{2 0 0}$ | $\mathbf{3 0 0}$ | $\mathbf{5 0 0}$ | $\mathbf{7 5 0}$ | $\mathbf{1 0 0 0}$ |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| Mean | 5930,769 | 5774,545 | 5764,78 | 5598,19 | 5628,787 | 5594,111 |
| Std Dev | 311,911 | 283,477 | 193,742 | 182,7593 | 107,9242 | 140,2526 |
| CI | 223,1277 | 202,7872 | 138,5947 | 130,7382 | 77,20433 | 100,3307 |
| CI/mean | $3,76 \%$ | $3,51 \%$ | $2,40 \%$ | $2,34 \%$ | $1,37 \%$ | $1,79 \%$ |

## Combination set of:

| s: | $\{11000$, | 12000 | 15200 | $11000\}$ |
| :---: | :---: | :---: | :---: | :---: |
| Q: | $\{400$, | 400 | 500 | $800\}$ |




## Statistical Information

|  | $\mathbf{1 0 0}$ | $\mathbf{2 0 0}$ | $\mathbf{3 0 0}$ | $\mathbf{5 0 0}$ | $\mathbf{7 5 0}$ | $\mathbf{1 0 0 0}$ |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| Mean | 3587,705 | 3204,043 | 3065,821 | 2943,074 | 2904,573 | 2884,157 |
| Std Dev | 79,67111 | 40,15379 | 46,1676 | 34,69438 | 18,22601 | 28,82404 |
| CI | 56,99328 | 28,72429 | 33,02631 | 24,81886 | 13,0381 | 20,61948 |
| Cl/mean | $1,59 \%$ | $0,90 \%$ | $1,08 \%$ | $0,84 \%$ | $0,45 \%$ | $0,71 \%$ |

## Sensitivity analysis: Tabu search parameters

For the table of the iterations, the values are in the following tables.

$$
11.2(\mathrm{~s}, \mathrm{Q}) \quad \Delta s=100, \Delta Q=5
$$

D.f.: 1
Initial
1600
$866\}$

| $\mathrm{T}^{*}(\mathrm{~s})$ | $\mathrm{F}^{*}$ (m.u.) |
| ---: | ---: |
| 0 | 3203.5 |
| 31.159 | 3172.17 |
| 46.668 | 3107.66 |
| 65.978 | 3068.93 |
| 72.105 | 2999.57 |
| 97.125 | 2995.18 |
| 117.125 | 2969.78 |
| 137.72 | 2897.14 |
| 160.799 | 2848.64 |
| 181.156 | 2846.77 |
| 249.878 | 2828.62 |
| 267.385 | 2808.41 |
| 284.194 | 2805.36 |
| 301.334 | 2754.5 |
| 340.344 | 2699.18 |
| 360.779 | 2697.62 |
| 373.769 | 2673.77 |
| 482.19 | 2661.03 |
| 503.235 | 2657.3 |
| 523.702 | 2645.77 |
| 537.041 | 2619.41 |
| 546.961 | 2608.74 |
| 560.284 | 2581.41 |
| 621.237 | 2557.89 |
| 650.976 | 2545.79 |
| 676.326 | 2521.56 |
| 1389.07 | 2516.1 |
| 1403.38 | 2508.72 |
| 1461.7 | 2494.69 |
| 1859.29 | 2485.82 |
| 1995.99 | 2483 |
| 2474.28 | 2470.69 |
| 2931.11 | 2466.59 |
| 3600 | 2466.59 |
|  |  |


| D.f.: 0.95 <br> Initial Solution: \{10000, 13000, 16000, 13000, 433, 500, 612, 866\} |  |
| :---: | :---: |
| T* | F* |
| 0 | 3226.51 |
| 13.238 | 3152.2 |
| 24.54 | 3101.84 |
| 36.221 | 3035.42 |
| 51.22 | 2987.03 |
| 68.527 | 2970 |
| 85.648 | 2904.27 |
| 107.651 | 2877.45 |
| 125.039 | 2831.71 |
| 199.371 | 2802.4 |
| 217.277 | 2788.47 |
| 253.751 | 2784.14 |
| 283.184 | 2767.93 |
| 299.711 | 2732.94 |
| 317.084 | 2707.96 |
| 335.085 | 2685.51 |
| 353.222 | 2644.27 |
| 373.231 | 2640.43 |
| 393.999 | 2630.46 |
| 411.463 | 2609.1 |
| 427.527 | 2597.29 |
| 444.616 | 2587.19 |
| 513.11 | 2566.14 |
| 534.388 | 2560.7 |
| 566.134 | 2556.58 |
| 665.792 | 2554.93 |
| 694.656 | 2545.74 |
| 717.397 | 2532.58 |
| 734.938 | 2531.75 |
| 756.74 | 2525.8 |
| 797.979 | 2516.83 |
| 865.911 | 2493.48 |
| 948.527 | 2472.97 |
| 966.98 | 2470.91 |
| 2744.15 | 2469.58 |


| $\begin{aligned} & \text { D.f.: 0.9 } \\ & \text { Initial Solution: }\{10000,13000 \text {, } \\ & 16000,13000,433,500,612, \\ & 866\} \end{aligned}$ |  |
| :---: | :---: |
| T* | F* |
| 0 | 3290.79 |
| 19.109 | 3207.02 |
| 43.443 | 3158.33 |
| 62.841 | 3090.63 |
| 82.889 | 3045.71 |
| 106.794 | 3003.99 |
| 132.35 | 2965.65 |
| 161.856 | 2955.05 |
| 183.382 | 2878.9 |
| 209.762 | 2861.61 |
| 223.491 | 2833.21 |
| 237.736 | 2817.01 |
| 251.383 | 2792.1 |
| 268.481 | 2761.11 |
| 293.901 | 2757.33 |
| 310.564 | 2732.25 |
| 323.661 | 2707.98 |
| 336.575 | 2687.7 |
| 356.349 | 2664.11 |
| 369.904 | 2638.55 |
| 420.068 | 2634.82 |
| 434.925 | 2622.66 |
| 448.491 | 2609.12 |
| 482.871 | 2597.57 |
| 565.289 | 2589.84 |
| 587.389 | 2559.54 |
| 635.802 | 2550.2 |
| 741.329 | 2543.62 |
| 751.023 | 2525.85 |
| 772.514 | 2493.08 |
| 781.653 | 2469.34 |
| 1021.02 | 2462.19 |
| 1134.39 | 2461.95 |
| 1480.55 | 2456.85 |
| 3280.35 | 2444.14 |


| 3350.35 | 2467.57 | 3600 | 2444.14 |
| ---: | ---: | ---: | ---: | ---: |
| 3504.65 | 2458.83 |  |  |
| 3600 | 2458.83 |  |  |

D.f.:
Initial
16000
$866\}$

Initial Solution: \{10000, 13000, 16000, 13000, 433, 500, 612,

| $\mathrm{T}^{*}(\mathrm{~s})$ | $\mathrm{F}^{*}$ (m.u.) |
| ---: | ---: |
| 0 | 3219.94 |
| 49.689 | 3149.96 |
| 68.11 | 3113.3 |
| 94.26 | 3042.68 |
| 124.388 | 3027.39 |
| 142.923 | 2980.18 |
| 160.469 | 2899.97 |
| 186.094 | 2868.16 |
| 212.721 | 2828.63 |
| 241.312 | 2814.66 |
| 259.541 | 2762.56 |
| 276.657 | 2705.48 |
| 325.065 | 2685.57 |
| 339.122 | 2662.05 |
| 353.763 | 2655.21 |
| 371.188 | 2646.38 |
| 409.619 | 2637.95 |
| 422.849 | 2619.02 |
| 435.813 | 2610.4 |
| 628.769 | 2606.96 |
| 647.447 | 2572.61 |
| 723.096 | 2571.98 |
| 742.157 | 2554.76 |
| 761.242 | 2513.08 |
| 1017.22 | 2512.66 |
| 1128.19 | 2503.74 |
| 1153.53 | 2485.81 |
| 1737.78 | 2485.41 |
| 1832.26 | 2480.91 |
| 1869.06 | 2475.11 |
| 3600 | 2475.11 |
|  |  |


| D.f.: 0.8 <br> Initial Solution: \{10000, 13000, 16000, 13000, 433, 500, 612, 866\} |  |
| :---: | :---: |
| T* | F* |
| 0 | 3204.73 |
| 20.387 | 3156.92 |
| 44.216 | 3135.96 |
| 66.397 | 3056.15 |
| 85.29 | 3026.58 |
| 107.281 | 2979.71 |
| 128.544 | 2931.96 |
| 144.808 | 2896.44 |
| 164.103 | 2861.54 |
| 185.758 | 2849.75 |
| 208.193 | 2822.93 |
| 230.445 | 2809.52 |
| 265.557 | 2760.23 |
| 282.382 | 2736.34 |
| 297.541 | 2734.46 |
| 315.367 | 2708.38 |
| 332.118 | 2689.27 |
| 340.96 | 2682.73 |
| 350.147 | 2651.11 |
| 362.061 | 2643.47 |
| 378.163 | 2623.75 |
| 408.501 | 2611.86 |
| 451.951 | 2608.13 |
| 472.268 | 2585.78 |
| 515.414 | 2555.32 |
| 532.538 | 2555.14 |
| 555.033 | 2552.76 |
| 704.301 | 2535.18 |
| 727.871 | 2534.41 |
| 746.376 | 2509.06 |
| 787.278 | 2490.03 |
| 822.867 | 2488.05 |
| 1704.06 | 2487.9 |
| 2427.3 | 2485.14 |
| 2776.31 | 2484.11 |
| 3179.45 | 2484.09 |

D.f.: 0.8
Initial Solution: $\{10000,13000$, 16000, 13000, 433, 500, 612, 866\}

| ```D.f.: 0.75 Initial Solution: {10000, 13000, 16000, 13000, 433, 500, 612, 866}``` |  |
| :---: | :---: |
| T* | F* |
| 0 | 3181.22 |
| 45.617 | 3132.96 |
| 70.057 | 3077.65 |
| 92.245 | 3053.43 |
| 113.87 | 2971.19 |
| 140.757 | 2923.96 |
| 163.056 | 2880.31 |
| 180.021 | 2838.8 |
| 207.073 | 2821.75 |
| 237.42 | 2811.69 |
| 250.77 | 2803.93 |
| 272.727 | 2784.85 |
| 285.162 | 2743.9 |
| 302.569 | 2738.33 |
| 322.635 | 2711.24 |
| 339.079 | 2647.1 |
| 355.857 | 2640.45 |
| 371.839 | 2634.27 |
| 380.932 | 2633.46 |
| 402.372 | 2620.22 |
| 415.141 | 2590.88 |
| 449.63 | 2561.14 |
| 564.882 | 2546.91 |
| 630.716 | 2534.38 |
| 648.683 | 2517.26 |
| 1139.8 | 2516.39 |
| 1158.3 | 2502.65 |
| 1190.69 | 2482.8 |
| 1352.7 | 2464.46 |
| 2933.64 | 2459.56 |
| 3491.27 | 2451.64 |
| 3600 | 2451.64 |

D.f.: 0.75
Initial Solution: $\{10000,13000$, 16000, 13000, 433, 500, 612, 866\}

| $\begin{aligned} & \text { D.f.: } 0.7 \\ & \text { Initial Solution: }\{10000,13000 \text {, } \\ & 16000,13000,433,500,612,866\} \end{aligned}$ |  |
| :---: | :---: |
|  |  |
|  |  |
| $\mathrm{T}^{*}$ | F* |
| 0 | 3322.55 |
| 27.449 | 3179.74 |
| 45.831 | 3139.19 |
| 60.168 | 3049.7 |
| 86.976 | 3015.99 |
| 113.388 | 2974.4 |
| 140.53 | 2946.09 |
| 166.464 | 2883.42 |
| 202.778 | 2876.89 |
| 264.181 | 2864.05 |
| 282.575 | 2855.47 |
| 307.883 | 2829.09 |
| 333.619 | 2773.78 |
| 379.747 | 2749.18 |
| 393.007 | 2714.32 |
| 414.814 | 2699.37 |
| 430.742 | 2697.17 |
| 443.902 | 2673.22 |
| 456.624 | 2650.21 |
| 474.242 | 2594.82 |
| 521.451 | 2583.73 |
| 534.574 | 2570.17 |
| 551.983 | 2553.36 |
| 569.884 | 2541.33 |
| 603.785 | 2517.11 |
| 902.564 | 2513.22 |
| 1094.45 | 2502.15 |
| 1514.72 | 2479.21 |
| 2001.27 | 2468.25 |
| 2053.84 | 2461.25 |
| 2076.95 | 2459.14 |
| 2219.78 | 2454.36 |
| 2553.07 | 2449.63 |
| 3269.48 | 2442.02 |
| 3462.2 | 2435.01 |
| 3600 | 2435.01 |

$$
\Delta s=500, \Delta Q=25
$$

| D.f.: 1 <br> Initial Solution: \{10000, 13000, 16000, 13000, 433, 500, 612, 866\} |  | $\begin{aligned} & \text { D.f.: } 0.95 \\ & \text { Initial Solution: }\{10000,13000 \text {, } \\ & 16000,13000,433,500,612 \text {, } \\ & 866\} \end{aligned}$ |  | $\begin{aligned} & \text { D.f.: } 0.9 \\ & \text { Initial Solution: }\{10000,13000 \text {, } \\ & 16000,13000,433,500,612, \\ & 866\} \end{aligned}$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{T}^{*}(\mathrm{~s})$ | F*(m.u.) | T* | F* | T* | F* |
| 0 | 3225.46 | 0 | 3225.71 | 0 | 3192.12 |
| 8.402 | 2994.3 | 11.322 | 2961.3 | 9.762 | 3061.02 |
| 20.434 | 2805.09 | 25.504 | 2786.61 | 26.588 | 2765.09 |
| 33.55 | 2688.04 | 41.575 | 2677.76 | 41.655 | 2687.6 |
| 48.541 | 2586.02 | 58.245 | 2593.45 | 56.059 | 2579.75 |
| 75.759 | 2548.67 | 78.873 | 2536.77 | 72.197 | 2570.96 |
| 100.59 | 2541.92 | 151.964 | 2530.64 | 92.677 | 2554.95 |
| 121.136 | 2506.08 | 243.173 | 2521.88 | 129.344 | 2552.65 |
| 866.842 | 2501.5 | 605.698 | 2511.32 | 145.478 | 2536.6 |
| 946.213 | 2497.92 | 3600 | 2511.32 | 344.877 | 2522.08 |
| 3600 | 2497.92 |  |  | 3600 | 2522.08 |


| D.f.: 0.85 <br> Initial Solution: \{10000, 13000, 16000, 13000, 433, 500, 612, 866\} |  |
| :---: | :---: |
| T* (s) | F*(m.u.) |
| 0 | 3288.4 |
| 11.252 | 2966.59 |
| 24.252 | 2781.09 |
| 43.025 | 2665.22 |
| 63.001 | 2649.19 |
| 92.081 | 2605.98 |
| 121.538 | 2599.45 |
| 167.379 | 2597.23 |
| 188.044 | 2555.87 |
| 213.255 | 2507.42 |
| 405.734 | 2495.83 |
| 1762.59 | 2491.59 |
| 1904.99 | 2482.76 |
| 3600 | 2482.76 |


| D.f.: 0.8 <br> Initial Solution: \{10000, 13000, 16000, 13000, 433, 500, 612, 866\} |  |
| :---: | :---: |
| $T^{*}$ | F* |
| 0 | 3242.45 |
| 16.526 | 3041.52 |
| 34.663 | 2781.01 |
| 54.595 | 2660.17 |
| 75.055 | 2590.94 |
| 104.482 | 2574.41 |
| 125.112 | 2544.64 |
| 145.646 | 2543.26 |
| 170.714 | 2536.44 |
| 258.135 | 2523.81 |
| 274.485 | 2510.61 |
| 299.968 | 2502.3 |
| 503.399 | 2476.05 |
| 3600 | 2476.05 |


| $\begin{aligned} & \text { D.f.: } 0.75 \\ & \text { Initial Solution: }\{10000,13000 \text {, } \\ & 16000,13000,433,500,612, \\ & 866\} \end{aligned}$ |  |
| :---: | :---: |
| $T^{*}$ | F* |
| 0 | 3231.51 |
| 19.966 | 2967.73 |
| 44.636 | 2762.35 |
| 65.172 | 2692.44 |
| 85.832 | 2586.16 |
| 122.764 | 2557.48 |
| 173.249 | 2522.33 |
| 533.19 | 2505.06 |
| 706.38 | 2493.65 |
| 3600 | 2493.65 |

## D.f.: 0.7

Initial Solution: $\{10000,13000,16000,13000,433$,
$500,612,866\}$

| $\mathrm{T}^{*}$ | $\mathrm{~F}^{*}$ |
| ---: | ---: |
| 0 | 3282.55 |
| 20.238 | 2978.13 |
| 36.178 | 2792.99 |
| 56.634 | 2734.14 |
| 81.644 | 2635 |
| 118.376 | 2631.46 |
| 138.927 | 2550.21 |
| 346.996 | 2542.11 |
| 379.4 | 2530.43 |
| 490.551 | 2530.04 |
| 684.351 | 2520.13 |
| 3600 | 2520.13 |

$$
\Delta s=1000, \quad \Delta Q=50
$$

| ```D.f.: } Initial Solution: {10000, 13000, 16000, 13000, 433, 500, 612, 866}``` |  | D.f.: 0.95 <br> Initial Solution: \{10000, 13000, 16000, 13000, 433, 500, 612, 866\} |  | ```D.f.: 0.9 Initial Solution: {10000, 13000, 16000, 13000, 433, 500, 612, 866}``` |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| T* $(\mathrm{s})$ | $\mathrm{F}^{*}$ (m.u.) | $\mathrm{T}^{*}(\mathrm{~s})$ | $\mathrm{F}^{*}$ (m.u.) | T* s ) | F*(m.u.) |
| 0 | 3181.32 | 0 | 3274.07 | 0 | 3248.45 |
| 8.64 | 2888.32 | 8.688 | 2861.44 | 9.865 | 2853.32 |
| 17.405 | 2678.84 | 19.24 | 2704.57 | 27.581 | 2688.33 |
| 32.386 | 2626.54 | 37.375 | 2668.45 | 46.138 | 2677.1 |
| 59.699 | 2624.62 | 51.812 | 2661.94 | 64.411 | 2617.79 |
| 101.511 | 2578.95 | 60.873 | 2608.86 | 116.358 | 2611.33 |
| 114.837 | 2567.18 | 82.267 | 2601.83 | 129.286 | 2609.11 |
| 138.131 | 2562.15 | 105.57 | 2584.42 | 178.649 | 2606.74 |
| 176.276 | 2539.19 | 245.903 | 2546.14 | 193.171 | 2601.13 |
| 3600 | 2539.19 | 3600 | 2546.14 | 285.985 | 2569.92 |
|  |  |  |  | 3210.42 | 2551.15 |
|  |  |  |  | 3233.58 | 2544.8 |
|  |  |  |  | 3323.13 | 2540.81 |
|  |  |  |  | 3600 | 2540.81 |


| D.f.: 0.85 <br> Initial Solution 16000, 13000, 866\} <br> $\mathrm{T}^{*}$ (s) | $\begin{aligned} & \text { 0000, 13000, } \\ & 3,500,612, \\ & \text { F*(m.u. }^{*} \end{aligned}$ | D.f.: 0.8 <br> Initial Solution: <br> 16000, 13000, <br> 866\} <br> $\mathrm{T}^{*}$ (s) | $\begin{aligned} & \text { 10000, 13000, } \\ & 3,500,612, \\ & \text { F*(m.u.) }^{*} \end{aligned}$ | D.f.: 0.75 <br> Initial Solution 16000, 13000, 866\} <br> T* (s) | $\begin{gathered} \text { 0000, 13000, } \\ 3,500,612, \\ \text { F*(m.u. }^{*} \end{gathered}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 3300.09 | 0 | 3258.73 | 0 | 3214.53 |
| 14.213 | 2848.22 | 17.067 | 2860.73 | 14.261 | 2884.1 |
| 37.971 | 2820.47 | 36.16 | 2692.01 | 33.02 | 2671.71 |
| 59.694 | 2787.72 | 54.697 | 2666.93 | 51.134 | 2658.8 |
| 78.167 | 2598.21 | 71.242 | 2603.5 | 64.709 | 2652.71 |
| 339.769 | 2585.8 | 172.21 | 2583.05 | 79.802 | 2644.71 |
| 688.893 | 2573.77 | 194.781 | 2571.09 | 98.143 | 2634.91 |
| 748.324 | 2554.05 | 209.191 | 2569.92 | 125.843 | 2588.56 |
| 3600 | 2554.05 | 459.081 | 2565.51 | 135.14 | 2582.55 |
|  |  | 2272.06 | 2548.17 | 228.53 | 2553.77 |
|  |  | 2868.67 | 2546.71 | 279.972 | 2529.5 |
|  |  | 3600 | 2546.71 | 3600 | 2529.5 |


| D.f.: 0.7 <br> Initial Solution: $13000,433,500,$ |  |
| :---: | :---: |
| T* (s) |  |
| 0 | 3206,32 |
| 19,809 | 2895,86 |
| 34,172 | 2694,42 |
| 56,515 | 2669,22 |
| 74,949 | 2648,37 |
| 107,783 | 2643,84 |
| 124,677 | 2642,72 |
| 133,757 | 2622,32 |
| 147,863 | 2580,12 |
| 170,329 | 2565,35 |
| 216,528 | 2559,83 |
| 265,346 | 2540,61 |
| 3600 | 2540,61 |

## 11. Appendix A: OUtput data



Figure 11.1: Iteration graph of the combinations of D.f. with $\Delta s=100, \Delta Q=5$


Figure 11.2: Iteration graph of the combinations of D.f. with $\Delta s=500, \Delta Q=25$


Figure 11.3: Iteration graph of the combinations of D.f. with $\Delta s=1000, \Delta Q=50$
11.2.2 ( $\mathbf{s}, \mathbf{S}$ )

$$
\Delta s=100, \quad \Delta S=100
$$

| D.f.: 1 |  | D.f.: 0.95 |  | D.f.: 0.9 |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| T* (s) | F*(m.u.) | T* (s) | F*(m.u.) | T* $(\mathrm{s})$ | $\mathrm{F}^{*}$ (m.u.) |
| 11.463 | 4170.5 | 12.23 | 4179.74 | 13.391 | 4158.95 |
| 22.552 | 4165.24 | 22.393 | 4105.23 | 31.76 | 4152.44 |
| 37.151 | 4151.17 | 56.274 | 4051.97 | 44.246 | 4140.76 |
| 56.328 | 4121.96 | 70.124 | 4051.27 | 63.099 | 4103.4 |
| 100.294 | 4119.14 | 76.829 | 3947.27 | 81.826 | 4100.16 |
| 120.518 | 4100.26 | 107.607 | 3937.64 | 95.428 | 4067.52 |
| 134.086 | 4080.71 | 114.118 | 3914.82 | 121.491 | 4056.02 |
| 145.121 | 4069.27 | 120.538 | 3889.69 | 137.876 | 4014.08 |
| 156.619 | 4026.14 | 130.442 | 3870.8 | 152.253 | 4008.19 |
| 167.421 | 4014.15 | 136.913 | 3824.68 | 164.004 | 4006.89 |
| 176.533 | 4012.72 | 150.751 | 3822.94 | 176.02 | 3962.83 |
| 191.861 | 3980.66 | 162.3 | 3776.83 | 191.372 | 3953.5 |
| 206.54 | 3974.62 | 202.504 | 3737.76 | 202.025 | 3914.03 |
| 231.721 | 3971.99 | 234.131 | 3720.73 | 259.839 | 3851.74 |
| 243.112 | 3952.17 | 250.377 | 3706.79 | 274.601 | 3848.42 |
| 262.683 | 3923.61 | 294.563 | 3677.83 | 285.572 | 3819.31 |
| 279.576 | 3843.08 | 304.48 | 3648.89 | 300.949 | 3800.92 |
| 310.944 | 3812.98 | 314.125 | 3641.44 | 307.587 | 3794.98 |
| 320.829 | 3805.49 | 325.59 | 3616.27 | 314.079 | 3705.58 |
| 334.313 | 3752.68 | 335.312 | 3608.78 | 363.924 | 3688.34 |
| 343.356 | 3749.94 | 349.527 | 3593.44 | 375.444 | 3682.07 |
| 350.407 | 3720.88 | 356.171 | 3592.36 | 390.834 | 3657.68 |
| 357.021 | 3715.07 | 371.945 | 3583.23 | 418.529 | 3628.06 |
| 366.248 | 3635.19 | 383.198 | 3553.08 | 433.513 | 3605.36 |
| 401.985 | 3597.56 | 398.721 | 3549.68 | 446.94 | 3601.83 |
| 413.152 | 3587.79 | 411.316 | 3542.08 | 481.099 | 3595.67 |
| 422.518 | 3546.77 | 427.119 | 3525.64 | 496.789 | 3558.58 |
| 455.134 | 3524.5 | 489.006 | 3510.15 | 523.643 | 3522.87 |
| 463.653 | 3501.02 | 515.906 | 3497.55 | 538.98 | 3480.06 |
| 477.715 | 3489.58 | 526.437 | 3472.62 | 587.719 | 3422.02 |
| 484.895 | 3484.26 | 553.107 | 3460.54 | 604.737 | 3381.07 |
| 517.637 | 3439.39 | 1538.14 | 3460.35 | 649.957 | 3374.02 |
| 528.086 | 3423.48 | 1590.47 | 3446.58 | 713.756 | 3371.97 |
| 543.276 | 3323.93 | 2166.4 | 3439.63 | 824.32 | 3370.73 |
| 652.457 | 3315.01 | 2380.35 | 3430.92 | 848.592 | 3367.12 |
| 711.039 | 3302.96 | 2401.07 | 3419.54 | 935.186 | 3366.27 |
| 759.36 | 3291.57 | 2428.35 | 3413.62 | 1050.07 | 3337.42 |
| 850.724 | 3278.6 | 2506.18 | 3398.91 | 1084.67 | 3324.35 |
| 1238.8 | 3278.48 | 3241.72 | 3397.27 | 1209.31 | 3323.49 |
| 1316.66 | 3271.22 | 3387.26 | 3383.61 | 2117.8 | 3312.93 |
| 1781.69 | 3257.36 | 3604.96 | 3378.33 | 2169.33 | 3312.34 |


| 1835.67 | 3254.91 | 4945 | 3375.06 | 2428.67 | 3304.86 |
| ---: | ---: | ---: | ---: | ---: | ---: |
| 1852.64 | 3248.72 | 5242.59 | 3362.18 | 2667.32 | 3299.79 |
| 2113.31 | 3248.49 | 5562.76 | 3354.66 | 2747.04 | 3270.64 |
| 2569 | 3237.66 | 5770.62 | 3345.78 | 2782.81 | 3265.03 |
| 3401.18 | 3234.39 | 6628.62 | 3323.82 | 2799.4 | 3248.34 |
| 4566.82 | 3229.12 | 6636.52 | 3322.86 | 3449.04 | 3247.98 |
| 7200 | 3229.12 | 6662.46 | 3311.98 | 3559.67 | 3247.24 |
|  |  | 7200 | 3311.98 | 3586.71 | 3231.6 |
|  |  |  |  | 5288.5 | 3226.83 |
|  |  |  | 6175.19 | 3210.94 |  |
|  |  |  | 7200 | 3210.94 |  |


| D.f.: 0.85 |  | D.f.: 0.8 |  | D.f.: 0.75 |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| T* s ) | F*(m.u.) | T* (s) | F*(m.u.) | T* s ) | F*(m.u.) |
| 15.292 | 4158.89 | 18.202 | 4161.25 | 18.079 | 4162.64 |
| 27.894 | 4155.96 | 28.967 | 4143.62 | 34.161 | 4138.08 |
| 43.858 | 4147.19 | 44.289 | 4113.18 | 58.502 | 4064.98 |
| 59.182 | 4124.09 | 79.719 | 4102.38 | 84.794 | 3999 |
| 70.658 | 4076.65 | 90.154 | 4093.84 | 99.041 | 3973.05 |
| 92.579 | 4047.57 | 104.545 | 4075.02 | 114.284 | 3958.01 |
| 105.758 | 4036.71 | 115.652 | 4039.26 | 129.873 | 3934.61 |
| 125.471 | 4029.34 | 126.219 | 4023.27 | 144.39 | 3917.24 |
| 131.594 | 3988 | 141.301 | 4009.16 | 155.628 | 3869.37 |
| 137.558 | 3979.85 | 152.297 | 3985.9 | 166.673 | 3867.12 |
| 147.713 | 3963.23 | 181.271 | 3939.34 | 182.972 | 3858.42 |
| 157.586 | 3927.96 | 192.665 | 3906.88 | 203.318 | 3823.07 |
| 168.463 | 3875.87 | 205.141 | 3903.43 | 303.093 | 3817.34 |
| 185.198 | 3859.65 | 211.612 | 3842.36 | 368.248 | 3795.3 |
| 201.86 | 3839.61 | 221.578 | 3830.28 | 388.453 | 3790.16 |
| 212.255 | 3820.97 | 265.542 | 3827.04 | 406.724 | 3745.6 |
| 223.024 | 3791.6 | 297.071 | 3805.97 | 450.309 | 3738.2 |
| 233.702 | 3789.19 | 312.094 | 3792.81 | 461.651 | 3712.45 |
| 243.638 | 3780.84 | 328.754 | 3782.19 | 473.268 | 3674.54 |
| 273.889 | 3772.72 | 343.39 | 3759.33 | 503.216 | 3648.24 |
| 283.358 | 3769.59 | 358.812 | 3753.36 | 523.278 | 3632.77 |
| 298.006 | 3747.83 | 373.195 | 3726.63 | 541.223 | 3623.02 |
| 308.769 | 3694.87 | 417.528 | 3716.31 | 555.826 | 3617.91 |
| 329.313 | 3690.37 | 432.771 | 3707.16 | 571.713 | 3613.08 |
| 339.58 | 3659.88 | 459.492 | 3677.21 | 587.122 | 3609.65 |
| 350.532 | 3638.99 | 470.866 | 3662.22 | 597.321 | 3566.58 |
| 372.825 | 3617.28 | 486.581 | 3647.75 | 608.124 | 3550.79 |
| 386.363 | 3591.41 | 497.584 | 3635.54 | 619.032 | 3549.7 |
| 397.727 | 3568.37 | 504.712 | 3588.74 | 644.596 | 3531.43 |
| 422.221 | 3565.74 | 535.074 | 3586.95 | 654.157 | 3530.48 |
| 435.973 | 3558.05 | 551.48 | 3565.45 | 676.339 | 3448.93 |


| 451.514 | 3543.37 | 574.368 | 3550.72 | 698.361 | 3433.22 |
| ---: | ---: | ---: | ---: | ---: | ---: |
| 476.318 | 3539.84 | 611.158 | 3530.4 | 713.323 | 3419.23 |
| 502.345 | 3536.97 | 646.819 | 3493.39 | 748.264 | 3396.84 |
| 517.433 | 3524.49 | 761.224 | 3475.52 | 764.762 | 3385.16 |
| 529.232 | 3514.44 | 799.697 | 3467.35 | 775.379 | 3372.21 |
| 539.457 | 3502.33 | 891.802 | 3467.05 | 801.779 | 3370.7 |
| 551.016 | 3491.39 | 968.173 | 3464.74 | 812.932 | 3364.62 |
| 589.426 | 3473.05 | 1064.33 | 3461.54 | 830.755 | 3351.87 |
| 601.821 | 3452.32 | 1075.44 | 3435.1 | 836.955 | 3347.11 |
| 636.326 | 3444.86 | 2057.48 | 3425.59 | 859.178 | 3345.7 |
| 683.181 | 3430.75 | 2074.15 | 3408.87 | 1033.31 | 3337.94 |
| 1672.28 | 3425.09 | 2136 | 3407.56 | 1054.66 | 3332.25 |
| 2391.36 | 3420.69 | 2171.95 | 3404.69 | 1082.09 | 3331.42 |
| 2468.84 | 3390.53 | 2878.21 | 3396.23 | 1337.99 | 3326.92 |
| 2813.79 | 3385.22 | 2967.18 | 3391.46 | 1382.38 | 3300.98 |
| 7200 | 3385.22 | 3736.91 | 3387.45 | 1426.16 | 3298.55 |
|  |  | 3773.08 | 3387.41 | 1459.88 | 3297.6 |
|  |  | 6066.07 | 3385.08 | 1732.09 | 3285.46 |
|  |  | 6084.49 | 3365.54 | 1767.05 | 3268.58 |
|  |  | 6350.15 | 3362.36 | 1840.71 | 3249.55 |
|  |  | 6372.56 | 3361.72 | 1889.02 | 3242.58 |
|  |  | 6394.84 | 3359.91 | 2091.14 | 3234.37 |
|  |  | 6415.59 | 3353.5 | 2107.64 | 3231.23 |
|  |  | 6422.37 | 3352.47 | 2370.11 | 3222.52 |
|  |  | 6449.52 | 3339.48 | 3889.64 | 3218.68 |
|  |  | 6821.48 | 3334.45 | 3954.17 | 3218.54 |
|  |  | 6862.97 | 3329.98 | 4296.51 | 3212.78 |
|  |  | 6875.49 | 3323.63 | 4647.76 | 3212.27 |
|  | 6883.11 | 3304.29 | 5762.4 | 3210.37 |  |
|  |  | 7172.53 | 3303.31 | 7200 | 3210.37 |
|  |  |  | 3303.31 |  |  |


| D.f.: $\mathbf{0 . 7}$ |  |
| ---: | :--- |
|  |  |
| $\mathrm{T}^{*}$ (s) | $\mathrm{F}^{*}$ (m.u.) |
| 19.041 | 4171.34 |
| 35.995 | 4171.21 |
| 52.008 | 4106.61 |
| 63.982 | 4066.87 |
| 74.202 | 4052.74 |
| 103.954 | 4044.24 |
| 119.665 | 4005.16 |
| 130.193 | 3969.69 |
| 140.825 | 3944.14 |
| 156.44 | 3930.42 |

11. Appendix A: Output data

| 168.092 | 3894.69 |
| ---: | ---: |
| 229.526 | 3873.4 |
| 249.939 | 3856.99 |
| 264.674 | 3835.08 |
| 276.88 | 3788.92 |
| 301.253 | 3740.61 |
| 327.588 | 3713.46 |
| 342.572 | 3706.67 |
| 373.613 | 3696.31 |
| 389.017 | 3659.79 |
| 419.866 | 3644.37 |
| 497.25 | 3640.94 |
| 518.348 | 3637.45 |
| 545.544 | 3610.49 |
| 571.544 | 3609.05 |
| 582.964 | 3603.73 |
| 604.506 | 3577.93 |
| 656.525 | 3555.55 |
| 671.795 | 3538.72 |
| 725.928 | 3536.83 |
| 755.346 | 3515.74 |
| 803.461 | 3471.81 |
| 874.204 | 3461.09 |
| 886.775 | 3451.33 |
| 1022.17 | 3444.94 |
| 2308.18 | 3444.53 |
| 2324.03 | 3439.02 |
| 2347.21 | 3434.77 |
| 2483.99 | 3429.69 |
| 2507.09 | 3427.05 |
| 2659.21 | 3404.79 |
| 2833.77 | 3401.63 |
| 5337.97 | 3394.29 |
| 7200 | 3394.29 |
|  |  |

$$
\Delta s=500, \quad \Delta S=500
$$

| D.f.: 1 |  | D.f.: 0.95 |  | D.f.: 0.9 |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| T* (s) | F*(m.u.) | T* (s) | F*(m.u.) | T* (s) | $\mathrm{F}^{*}$ (m.u.) |
| 9.859 | 4149.12 | 7.798 | 4178.89 | 7.869 | 4161.82 |
| 15.859 | 4105.31 | 11.734 | 4033.89 | 14.441 | 3990.14 |
| 24.469 | 4087.81 | 17.798 | 3983.6 | 28.169 | 3945.44 |
| 31.149 | 4008.51 | 22.255 | 3860.15 | 34.416 | 3909.25 |
| 42.174 | 3941.93 | 28.829 | 3825.77 | 38.596 | 3847.03 |
| 51.172 | 3939.51 | 35.602 | 3755.43 | 47.358 | 3826.48 |
| 55.532 | 3803.65 | 44.813 | 3733.96 | 51.462 | 3735.65 |
| 69.327 | 3760.69 | 49.299 | 3674.47 | 57.921 | 3720.02 |
| 187.335 | 3706.04 | 100.238 | 3667.83 | 106.366 | 3717.75 |
| 217.261 | 3695.22 | 127.223 | 3655.59 | 127.614 | 3698.34 |
| 259.2 | 3686.11 | 375.455 | 3655.34 | 134.028 | 3674.42 |
| 306.172 | 3681.07 | 806.72 | 3653.34 | 2886.58 | 3654.12 |
| 850.768 | 3676.55 | 1549.66 | 3649.84 | 2978.69 | 3650.22 |
| 865.592 | 3670.35 | 2297.54 | 3642.31 | 3068.74 | 3635.48 |
| 875.24 | 3667.05 | 5684.01 | 3638.03 | 4160.74 | 3630.01 |
| 2078.89 | 3664.66 | 6633.75 | 3601.4 | 4181.81 | 3624.98 |
| 2112.47 | 3655.59 | 7045.53 | 3580.94 | 4357.25 | 3619.39 |
| 7200 | 3655.59 | 7052.35 | 3567.57 | 4560.03 | 3604.89 |
|  |  | 7200 | 3567.57 | 4765.72 | 3596.33 |
|  |  |  |  | 4792.74 | 3594.83 |
|  |  |  |  | 4831.1 | 3587.35 |
|  |  |  |  | 4859.9 | 3562.17 |
|  |  |  |  | 4881.56 | 3534.03 |
|  |  |  |  | 5006.25 | 3521.13 |
|  |  |  |  | 5181.46 | 3507.3 |
|  |  |  |  | 5185.72 | 3507.18 |
|  |  |  |  | 5312.12 | 3499.09 |
|  |  |  |  | 5407.73 | 3496.75 |
|  |  |  |  | 5455.32 | 3482.33 |
|  |  |  |  | 7200 | 3482.33 |


| D.f.: 0.85 |  | D.f.: 0.8 |  | D.f.: 0.75 |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| T* (s) | $\mathrm{F}^{*}$ (m.u.) | T* (s) | $\mathrm{F}^{*}$ (m.u.) | T* (s) | F*(m.u.) |
| 11.6 | 4160.62 | 7.764 | 4175.17 | 9.786 | 4134.52 |
| 18.472 | 4004.82 | 11.744 | 4004.82 | 13.739 | 3980.06 |
| 27.34 | 3967.26 | 19.371 | 3998.05 | 19.344 | 3956.28 |
| 31.685 | 3928.79 | 27.136 | 3951.52 | 23.097 | 3867.32 |
| 41.943 | 3684.79 | 32.461 | 3945.36 | 30.714 | 3846.23 |
| 57.205 | 3637.66 | 36.024 | 3796.06 | 36.473 | 3778.85 |
| 69.806 | 3634.35 | 43.445 | 3769.54 | 41.839 | 3770.64 |
| 87.015 | 3623.29 | 49.079 | 3733.78 | 45.399 | 3693.04 |
| 93.146 | 3592.98 | 54.528 | 3720.14 | 50.853 | 3677.24 |
| 415.72 | 3590.12 | 119.3 | 3703.64 | 54.476 | 3648.69 |
| 422.18 | 3586.64 | 133.19 | 3694.48 | 1544.69 | 3646.78 |
| 438.53 | 3561.38 | 179.35 | 3679.44 | 1553.45 | 3645.45 |
| 703.04 | 3557.79 | 281.72 | 3652.42 | 1674.06 | 3636.12 |
| 709.8 | 3554.81 | 494.05 | 3646.5 | 1680.61 | 3621.75 |
| 775.48 | 3554.17 | 1516.3 | 3597.28 | 1734.73 | 3616.06 |
| 1709.1 | 3549.66 | 1846.9 | 3569 | 1748.08 | 3610.89 |
| 1985.9 | 3543.8 | 1863.5 | 3558.07 | 1769.41 | 3609.89 |
| 2033.2 | 3539.81 | 1877.5 | 3549.63 | 1772.98 | 3609.06 |
| 2041.5 | 3527.07 | 2002.2 | 3541.96 | 1851.89 | 3599.69 |
| 2356.9 | 3524.23 | 2010.4 | 3522.3 | 1906.95 | 3593.34 |
| 2563.1 | 3513.82 | 2155.6 | 3513.3 | 2005.33 | 3567.68 |
| 2753.4 | 3509.52 | 2387.3 | 3505.55 | 2083.63 | 3562.5 |
| 2991.6 | 3508.08 | 2687.5 | 3457.18 | 2109.9 | 3545.58 |
| 3052.7 | 3507.25 | 2703.1 | 3425.42 | 2179.84 | 3539.51 |
| 3123.3 | 3505.32 | 2744 | 3420.73 | 2223.14 | 3526.4 |
| 3309.8 | 3503.16 | 3323.4 | 3419.22 | 2352.76 | 3508.01 |
| 3516.1 | 3501.5 | 4240.5 | 3417.13 | 2394.38 | 3499.36 |
| 4494.6 | 3499.62 | 4387.9 | 3370.82 | 2549.81 | 3480.15 |
| 4849.1 | 3481.67 | 4396.3 | 3362.94 | 2961.44 | 3461.22 |
| 4857.5 | 3459.53 | 5972.4 | 3362.55 | 2991.2 | 3454.98 |
| 5037.9 | 3421.36 | 6247.6 | 3361.08 | 3068.14 | 3436.7 |
| 5153.8 | 3413.12 | 7200 | 3361.08 | 3076.98 | 3419.94 |
| 5409 | 3389.17 |  |  | 3247.38 | 3417.18 |
| 5438.1 | 3354.46 |  |  | 3294.3 | 3412 |
| 5455.9 | 3353.41 |  |  | 3305.06 | 3408.12 |
| 5562.1 | 3343.02 |  |  | 3317.92 | 3395.66 |
| 5628.9 | 3340.45 |  |  | 3337.15 | 3392.32 |
| 5702.9 | 3337.19 |  |  | 3625.56 | 3384.67 |
| 5729.7 | 3328.42 |  |  | 3632.38 | 3362.1 |
| 5777.4 | 3326 |  |  | 4024.27 | 3359.11 |
| 5797.1 | 3320.25 |  |  | 4136.14 | 3351.06 |
| 5803.6 | 3300.45 |  |  | 4176.66 | 3327.37 |
| 5820.5 | 3287.13 |  |  | 4184.37 | 3326.49 |
| 5895.7 | 3285.6 |  |  | 4201.77 | 3325.99 |


| 5966.6 | 3278.02 |  | 4242.94 | 3324.62 |
| :---: | :---: | :---: | :---: | :---: |
| 6053.7 | 3276.14 |  | 4254.12 | 3323.36 |
| 6185.9 | 3263.46 |  | 4353.57 | 3300.93 |
| 6205.8 | 3244.49 |  | 4438.04 | 3293.07 |
| 7021.4 | 3240.04 |  | 4494.2 | 3285 |
| 7028.5 | 3235.39 |  | 4556.52 | 3283.43 |
| 7046.1 | 3208.19 |  | 4640.32 | 3274.31 |
| 7200 | 3208.19 |  | 4668.3 | 3269.65 |
|  |  |  | 4893.66 | 3269.46 |
|  |  |  | 4976.25 | 3260.3 |
|  |  |  | 5476.48 | 3258.92 |
|  |  |  | 5628.65 | 3254.05 |
|  |  |  | 5711.77 | 3249.78 |
|  |  |  | 5909.28 | 3248.7 |
|  |  |  | 5929.37 | 3244.36 |
|  |  |  | 6009.76 | 3241.15 |
|  |  |  | 6200.74 | 3240.9 |
|  |  |  | 6392.27 | 3237.26 |
|  |  |  | 6576.96 | 3231.75 |
|  |  |  | 6603.49 | 3225.61 |
|  |  |  | 7200 | 3225.61 |


| D.f.: 0.7 |  |
| ---: | ---: |
| $\mathrm{~T}^{*}(\mathrm{~s})$ | $\mathrm{F}^{*}$ (m.u.) |
| 9.896 | 4140.69 |
| 13.806 | 4028.73 |
| 21.352 | 3985.26 |
| 29.083 | 3951.5 |
| 34.461 | 3921.62 |
| 37.979 | 3771.78 |
| 45.38 | 3755.14 |
| 48.909 | 3718.67 |
| 56.447 | 3714.95 |
| 75.159 | 3692.74 |
| 382.699 | 3688.06 |
| 390.483 | 3679.22 |
| 404.944 | 3678.21 |
| 418.08 | 3671.84 |
| 431.002 | 3662.09 |
| 2253.65 | 3637.22 |
| 2270.4 | 3636.95 |
| 2278.8 | 3632 |
| 2588.15 | 3631.18 |
| 2606.3 | 3597.84 |
| 2727.38 | 3595.9 |


| 2735.52 | 3585.2 |
| ---: | ---: |
| 2814.21 | 3580.96 |
| 2830.79 | 3576.71 |
| 2874.5 | 3567.64 |
| 2883.51 | 3552.72 |
| 2896.44 | 3545.99 |
| 2973.81 | 3541.87 |
| 3004.24 | 3521.06 |
| 3017.48 | 3509.52 |
| 3579.83 | 3499.99 |
| 3658.72 | 3495.1 |
| 3839.86 | 3492.4 |
| 3904.47 | 3481.45 |
| 3914.98 | 3472.45 |
| 3930.24 | 3470.87 |
| 3967.04 | 3465.55 |
| 3978.53 | 3459.3 |
| 3994.61 | 3443.08 |
| 4437.93 | 3440.75 |
| 4463.35 | 3435.16 |
| 4508.96 | 3430.61 |
| 4599.17 | 3424.36 |
| 5629.59 | 3421.03 |
| 5672.28 | 3420.49 |
| 5766.47 | 3415.86 |
| 5889.27 | 3415.67 |
| 5927.5 | 3409.01 |
| 7200 | 3409.01 |

$$
\Delta s=1000, \quad \Delta S=1000
$$

| D.f.: $\mathbf{1}$ |  | D.f.: $\mathbf{0 . 9 5}$ | D.f.: $\mathbf{0 . 9}$ |  |  |
| ---: | ---: | ---: | ---: | ---: | ---: |
| $\mathrm{T}^{*}(\mathrm{~s})$ | $\mathrm{F}^{*}$ (m.u.) | $\mathrm{T}^{*}(\mathrm{~s})$ | $\mathrm{F}^{*}(\mathrm{~m} . \mathrm{u})$. | $\mathrm{T}^{*}(\mathrm{~s})$ | $\mathrm{F}^{*}(\mathrm{~m} . \mathrm{u})$. |
| 8.898 | 4138.16 | 12.715 | 4136.79 | 14.602 | 4132.65 |
| 17.267 | 3982.32 | 22.66 | 4003.95 | 26.344 | 3977.81 |
| 39.61 | 3744.27 | 33.814 | 3999.33 | 54.025 | 3728.92 |
| 69.678 | 3722.73 | 41.725 | 3712.13 | 129.919 | 3720.07 |
| 84.529 | 3719.96 | 223.483 | 3711.26 | 194.356 | 3707.23 |
| 144.835 | 3718.94 | 245.03 | 3708.09 | 318.988 | 3690.07 |
| 293.872 | 3718.85 | 412.278 | 3705.78 | 4623.13 | 3689.26 |
| 320.545 | 3718.77 | 907.668 | 3693.8 | 4642.94 | 3672.86 |
| 397.044 | 3702.86 | 1977.64 | 3692.89 | 7200 | 3672.86 |
| 591.065 | 3702.19 | 5636.19 | 3691.1 |  |  |
| 1089.44 | 3698.2 | 6806.67 | 3684.42 |  |  |
| 7200 | 3698.2 | 7200 | 3684.42 |  |  |


| D.f.: 0.85 |  | D.f.: 0.8 |  | D.f.: 0.75 |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| T* (s) | $F^{*}$ (m.u.) | T* (s) | F*(m.u.) | T* (s) | F*(m.u.) |
| 9.056 | 4154.59 | 14.121 | 4160.59 | 10.898 | 4134.81 |
| 19.989 | 3862.92 | 25.356 | 3827.37 | 22.917 | 3842.08 |
| 31.192 | 3813.82 | 36.984 | 3816.77 | 34.217 | 3838.86 |
| 39.716 | 3725.77 | 48.343 | 3715.72 | 42.07 | 3726.46 |
| 251.79 | 3709.99 | 126.895 | 3711.45 | 52.491 | 3720.17 |
| 634.836 | 3707.46 | 631.016 | 3710.88 | 174.579 | 3711.66 |
| 891.271 | 3705.36 | 727.787 | 3710.62 | 287.276 | 3705.5 |
| 1038.01 | 3698.68 | 778.982 | 3708.37 | 407.766 | 3704.65 |
| 1677.63 | 3695.57 | 1059.97 | 3706.85 | 972.686 | 3701.66 |
| 1721.68 | 3691.85 | 1264.3 | 3698.6 | 2867.83 | 3676.87 |
| 1824.49 | 3685.64 | 1680.37 | 3691.81 | 2934.08 | 3668.82 |
| 2702.21 | 3669.71 | 2284.92 | 3682.01 | 3031.2 | 3663.52 |
| 2770.21 | 3668.74 | 2296.83 | 3674.6 | 3385.33 | 3656.32 |
| 3070.29 | 3653.91 | 2396.58 | 3663.98 | 3401.57 | 3631.5 |
| 4778.9 | 3644.05 | 2857.07 | 3656.7 | 3428.24 | 3624.16 |
| 6338 | 3616.12 | 4015.96 | 3640.99 | 3446.21 | 3615.16 |
| 6384.54 | 3604.88 | 7200 | 3640.99 | 3463.95 | 3601.61 |
| 6605.2 | 3602.49 |  |  | 3732.71 | 3572.14 |
| 7146.28 | 3595.73 |  |  | 3777.98 | 3564.87 |
| 7178.22 | 3579.33 |  |  | 3909.87 | 3551.82 |
| 7200 | 3579.33 |  |  | 4250.51 | 3540.28 |
|  |  |  |  | 4266.24 | 3531.47 |
|  |  |  |  | 4734.9 | 3529.81 |
|  |  |  |  | 4779.94 | 3516.35 |
|  |  |  |  | 4793.55 | 3509.71 |
|  |  |  |  | 5676.62 | 3508.66 |
|  |  |  |  | 6353.05 | 3497.62 |
|  |  |  |  | 6887.45 | 3487.16 |
|  |  |  |  | 7200 | 3487.16 |


| D.f.: $\mathbf{0 . 7}$ |  |
| ---: | ---: |
| $\mathrm{T}^{*}(\mathrm{~s})$ | $\mathrm{F}^{*}$ (m.u.) |
| 15.589 | 4149.72 |
| 28.409 | 3980.73 |
| 51.875 | 3735.75 |
| 58.695 | 3716.68 |
| 196.345 | 3710.58 |
| 247.091 | 3708.98 |
| 479.411 | 3707.49 |
| 662.38 | 3703.01 |
| 711.696 | 3694.73 |
| 2113.64 | 3660.53 |
| 2124.48 | 3658.82 |

11. Appendix A: Output data

| 2274.12 | 3651.39 |
| ---: | ---: |
| 2421.22 | 3650.44 |
| 2934.54 | 3637.75 |
| 3703.64 | 3634.12 |
| 3752.85 | 3626.9 |
| 3847.85 | 3625.73 |
| 3865.51 | 3623.66 |
| 3893.52 | 3620.31 |
| 4161.34 | 3619.92 |
| 4263.46 | 3615.67 |
| 4689.5 | 3611.74 |
| 4797.65 | 3610.57 |
| 4983.4 | 3598.96 |
| 4997.91 | 3586.91 |
| 5086.07 | 3563.01 |
| 5367.73 | 3554.99 |
| 5407.65 | 3535.04 |
| 5745.3 | 3525.07 |
| 5776.69 | 3524.26 |
| 5821.77 | 3517.82 |
| 6573.75 | 3515.94 |
| 6790.36 | 3503.2 |
| 7069.7 | 3496.75 |
| 7181.02 | 3467.79 |
| 7200 | 3467.79 |
|  |  |



Figure 11.4: Iteration graph of the combinations of D.f. with $\Delta s=100, \Delta S=100$


Figure 11.5: Iteration graph of the combinations of D.f. with $\Delta s=500, \Delta S=500$


Figure 11.6: Iteration graph of the combinations of D.f. with $\Delta s=1000, \Delta S=1000$
11.3 DTS, STS and DLS sensitivity analysis

$$
11.3 .1 \quad(\mathbf{s}, \mathbf{Q})
$$

| Dynamic TS | Static TS | DLS | Dynamic TS | Static TS | DLS |
| ---: | ---: | ---: | ---: | ---: | ---: |
| $\mathbf{t}^{*}(\mathbf{s})$ | $\mathbf{t}^{*}$ | $\mathbf{t}^{*}$ | $\mathbf{F}^{*}$ | $\mathbf{F}^{*}$ | 2501.5 |
| 2323.9 | 204.53 | 1829.3 | 2511.6 | 2506.93 | 2750.1 |
| 81.8 | 127 | 2859.5 | 2500.0 | 2508.01 | 2568.5 |
| 709.3 | 537.7 | 435.0 | 2473.6 | 2508.18 | 2529.9 |
| 3531.3 | 617 | 1791.8 | 2475.6 | 2485.89 | 2786.9 |
| 422.0 | 210 | 1253.4 | 2445.7 | 2473.57 | 3158.6 |
| 1905.6 | 388 | 1357.7 | 2474.4 | 2503.52 | 2821.7 |
| 1330.3 | 467 | 313.8 | 2482.4 | 2503.77 | 2497.1 |
| 1253.7 | 171.48 | 0.0 | 2462.7 | 2462.71 | 3222.3 |
| 2361.1 | 341 | 1407.5 | 2473.3 | 2501.68 | 2492.3 |

### 11.3.2 (s,S)

| Dynamic TS | Static TS | DLS | Dynamic TS | Static TS | DLS |
| ---: | ---: | ---: | ---: | ---: | ---: |
| $\mathbf{\mathbf { t } ^ { * } ( \mathbf { s } )}$ | $\mathbf{t}^{*}$ | $\mathbf{t}^{*}$ | $\mathbf{F}^{*}$ | $\mathbf{F}^{*}$ | $\mathbf{F}^{*}$ |
| 6936.0 | 2819.2 | 1198.3 | 3264.0 | 3212.2 | 3643.2 |
| 6074.0 | 2551.3 | 2094.5 | 3224.5 | 3224.8 | 6820.4 |
| 5698.2 | 4004.3 | 260.9 | 3339.8 | 3272.4 | 3536.2 |
| 7154.5 | 6772.0 | 3490.5 | 3302.9 | 3364.3 | 5487.1 |
| 6252.1 | 3021.4 | 5334.8 | 3230.3 | 3367.2 | 4908.5 |
| 6829.0 | 6775.2 | 1011.8 | 3327.6 | 3232.8 | 3539.7 |
| 6977.1 | 6986.8 | 366.7 | 3268.4 | 3278.5 | 3603.6 |
| 4285.2 | 5731.8 | 584.5 | 3199.0 | 3239.4 | 3696.9 |
| 6190.6 | 6342.3 | 1279.0 | 3299.2 | 3323.1 | 4496.5 |
| 7105.6 | 6785.8 | 2632.2 | 3319.6 | 3313.4 | 4165.3 |

### 11.4 Detailed cost of the results

11.4.1 Cost of the best solution ( $\mathbf{s}, \mathbf{Q}$ )

Solution of the $(s, Q)$ case is found in the following table, the cost is the daily cost. The daily cost results from dividing the cost of all the days by the number of days.

|  |  | Retailer | Wholesaler | Distributor | Factory |
| :--- | ---: | ---: | ---: | ---: | ---: |
| Backorder cost | [m.u.] | 20,27 | 66,23 | 50,50 | 26,97 |
| Holding cost | [m.u.] | 946,94 | 356,08 | 300,45 | 152,54 |
| Launching cost | [m.u.] | 177,35 | 132,95 | 115,85 | 77,35 |

## 11. Appendix A: Output data

11.4.2 Cost of the best solution ( $\mathbf{s}, \mathbf{S}$ )

|  |  | Retailer | Wholesaler | Distributor | Factory |
| :--- | ---: | ---: | ---: | ---: | ---: |
| Backorder cost | [m.u.] | 28,22 | 184,66 | 52,09 | 92,01 |
| Holding cost | [m.u.] | 1014,26 | 407,63 | 575,42 | 376,11 |
| Launching cost | [m.u.] | 137,70 | 130,50 | 91,90 | 63,20 |

11.4.3 Cost for the several service levels ( $\mathbf{s}, \mathbf{Q}$ ) case

| SL=0.92 |  |  | SL=0.97 |  |  | SL=0.95 |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Cost | SL | Daily cost | Cost | SL | Daily cost | Cost | SL | Daily cost |
| 1161110 | 0.897 | 2322.22 | 1320640 | 0.976 | 2641.28 | 1233540 | 0.968 | 2467.08 |
| 1164340 | 0.906 | 2328.68 | 1293070 | 0.954 | 2586.14 | 1218310 | 0.951 | 2436.62 |
| 1180460 | 0.921 | 2360.92 | 1307650 | 0.959 | 2615.30 | 1189520 | 0.936 | 2379.04 |
| 1182770 | 0.883 | 2365.54 | 1345850 | 0.988 | 2691.70 | 1209520 | 0.951 | 2419.04 |
| 1186770 | 0.811 | 2373.54 | 1339050 | 0.981 | 2678.10 | 1209340 | 0.957 | 2418.68 |
| 1170200 | 0.913 | 2340.40 | 1317370 | 0.982 | 2634.74 | 1193330 | 0.942 | 2386.66 |
| 1163380 | 0.907 | 2326.76 | 1286130 | 0.929 | 2572.26 | 1220370 | 0.940 | 2440.74 |
| 1156710 | 0.907 | 2313.42 | 1300500 | 0.988 | 2601.00 | 1193690 | 0.955 | 2387.38 |
| 1175810 | 0.882 | 2351.62 | 1326770 | 0.976 | 2653.54 | 1206290 | 0.938 | 2412.58 |
| 1204550 | 0.901 | 2409.10 | 1312880 | 0.951 | 2625.76 | 1177530 | 0.945 | 2355.06 |
| AVG | 0.893 | 2349.220 | AVG | 0.968 | 2629.982 | AVG | 0.948 | 2410.288 |
| Std Dev | 0.030 | 27.570 | Std Dev | 0.018 | 36.273 | Std Dev | 0.010 | 31.793 |
| IC | 0.021 | 19.722 | IC | 0.013 | 25.948 | IC | 0.007 | 22.743 |
| IC Min | 0.871 | 2329.498 | IC Min | 0.955 | 2604.034 | IC Min | 0.942 | 2387.545 |
| IC Max | 0.914 | 2368.942 | IC Max | 0.982 | 2655.930 | IC Max | 0.955 | 2433.031 |

Distribution of the costs:

SL=0.92

|  |  | Retailer | Wholesaler | Distributor | Factory |
| :--- | :--- | ---: | ---: | ---: | ---: |
| Backorder cost | [m.u.] | 26.71 | 75.36 | 70.76 | 53.87 |
| Holding cost | [m.u.] | 856.79 | 334.37 | 253.60 | 129.98 |
| Launching cost | [m.u.] | 181.50 | 136.05 | 129.05 | 80.85 |

$\mathrm{SL}=0.98$

|  |  | Retailer | Wholesaler | Distributor | Factory |
| :--- | :--- | ---: | ---: | ---: | ---: |
| Backorder cost | [m.u.] | 9.28 | 123.33 | 51.68 | 24.85 |
| Holding cost | [m.u.] | 1142.63 | 266.40 | 271.52 | 174.49 |
| Launching cost | [m.u.] | 206.70 | 138.35 | 121.50 | 81.75 |

11.4.1 Cost for the several service levels (s.S) case

| SL=0.88 |  |  | SL=0.97 |  |  | SL=0.94 |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Cost | SL | Daily cost | Cost | SL | Daily cost | Cost | SL | Daily cost |
| 1567630 | 0.882 | 3135.26 | 1727060 | 0.981 | 3454.12 | 1620470 | 0.913 | 3240.94 |
| 1591030 | 0.888 | 3182.06 | 1702830 | 0.964 | 3405.66 | 1616120 | 0.960 | 3232.24 |
| 1569540 | 0.860 | 3139.08 | 1694430 | 0.960 | 3388.86 | 1640220 | 0.958 | 3280.44 |
| 1564230 | 0.892 | 3128.46 | 1719400 | 0.970 | 3438.80 | 1604380 | 0.952 | 3208.76 |
| 1552890 | 0.875 | 3105.78 | 1744500 | 0.957 | 3489.00 | 1621810 | 0.950 | 3243.62 |
| 1553010 | 0.890 | 3106.02 | 1719910 | 0.953 | 3439.82 | 1611990 | 0.935 | 3223.98 |
| 1558870 | 0.893 | 3117.74 | 1679340 | 0.970 | 3358.68 | 1573370 | 0.930 | 3146.74 |
| 1566300 | 0.908 | 3132.60 | 1725850 | 0.976 | 3451.70 | 1594970 | 0.928 | 3189.94 |
| 1594160 | 0.858 | 3188.32 | 1712360 | 0.962 | 3424.72 | 1608110 | 0.934 | 3216.22 |
| 1549150 | 0.854 | 3098.30 | 1708930 | 0.976 | 3417.86 | 1605010 | 0.934 | 3210.02 |
| AVG | 0.880 | 3133.36 | AVG | 0.967 | 3426.92 | AVG | 0.939 | 3219.29 |
| Std Dev | 0.017 | 29.04 | Std Dev | 0.009 | 34.87 | Std Dev | 0.014 | 33.65 |
| IC | 0.012 | 20.78 | IC | 0.006 | 24.94 | IC | 0.010 | 24.07 |
| IC Min | 0.868 | 3112.59 | IC Min | 0.961 | 3401.98 | IC Min | 0.929 | 3195.22 |
| IC Max | 0.892 | 3154.14 | IC Max | 0.973 | 3451.87 | IC Max | 0.949 | 3243.36 |

Distribution of the costs:

SL=0.88

|  |  | Retailer | Wholesaler | Distributor | Factory |
| :--- | :--- | :--- | ---: | ---: | ---: |
| Backorder cost | [m.u.] | 77.59 | 264.96 | 80.13 | 112.16 |
| Holding cost | [m.u.] | 913.12 | 392.88 | 549.67 | 398.63 |
| Launching cost | [m.u.] | 135.10 | 125.80 | 87.40 | 59.25 |

SL=0.97

|  |  | Retailer | Wholesaler | Distributor | Factory |
| :--- | :--- | ---: | ---: | ---: | ---: |
| Backorder cost | [m.u.] | 9.24 | 64.24 | 102.90 | 56.12 |
| Holding cost | [m.u.] | 1236.34 | 622.55 | 423.52 | 455.72 |
| Launching cost | [m.u.] | 137.90 | 131.80 | 92.45 | 63.40 |

## 11. Appendix A: Output data

### 11.5 Other solutions found

A different solution was found using the tabu search. Other solutions different than the best are found, but the differences in the values of the solution are small so they are not considered. But a solution with a remarkable difference is found, that leads to a similar cost and on-hand inventory but has a different distribution of costs. This is expected as the solution has 8 variables, so a lot of combinations can be made and the optimal solution does not have to be unique. This solution found has these values (Table 11.1):

| Table 11.1: Solution for the (s,Q) case II |  |  |  |  |  |
| ---: | ---: | ---: | ---: | ---: | ---: |
|  |  | Factory | Distributor | Wholesaler | Retailer |
| $\boldsymbol{s}_{\boldsymbol{i}}$ | [pieces] | 12000 | 15300 | 12200 | 10200 |
| $\boldsymbol{Q}_{\boldsymbol{i}}$ | [pieces] | 856 | 552 | 510 | 353 |

To compare this with the other solution a bar figure is made:


The two solutions, the one of the previous chapter and this new one have similar total cost. However, the length of the confidence interval is different, as the interval of the new solution includes the interval of the first one.

|  | First solution | New solution |
| :---: | :---: | :---: |
| Daily Total cost 95\% C.I. | $[2372.56,2402,41]$ | $[2365.03,2428.01]$ |
| Service Level 95\% C.I. | $[0.937,0.957]$ | $[0.926,0.952]$ |


[^0]:    ${ }^{1}$ XBeerGame, Université Laval

[^1]:    ${ }^{2}$ On-hand inventory: Physical inventory (stock) at the echelon
    ${ }^{3}$ Backlogged orders: Orders that have not been satisfied in this time step or in the previous ones

[^2]:    ${ }^{4}$ Inventory position: on-hand stock + standing orders - backlogged orders

[^3]:    ${ }^{5}$ t.s.u.: Time step units
    ${ }^{6}$ m.u.: Monetary units

[^4]:    ${ }^{7}$ FIFO: First in, first out.

[^5]:    ${ }^{8}$ The index 0 represents the manufacturer, who ships the material when it is ordered.

[^6]:    ${ }^{9}$ Greater figures can be found in 11.2
    ${ }^{10}$ D.f.: Decreasing Factor for the variation parameters

