

Solving electric market quadratic problems by Branch and Fix Coordination methods ^{*}

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Abstract. *The electric market regulation in Spain (MIBEL) establishes the rules for bilateral and futures contracts in the day-ahead optimal bid problem. Our model allows a price-taker generation company to decide the unit commitment of the thermal units, the economic dispatch of the bilateral and futures contracts between the thermal units and the optimal sale bids for the thermal units observing the MIBEL regulation. The uncertainty of the spot prices is represented through scenario sets. We solve this model on the framework of the Branch and Fix Coordination methodology as a quadratic two-stage stochastic problem. In order to gain computational efficiency, we use scenario clusters and propose to use perspective cuts. Numerical results are reported.*

Keywords: *Liberalized Electricity Market, Optimal Bid, Stochastic Programming, Quadratic Branch-and-Fix Coordination*

1 Introduction

This work is applied to the Iberian Electricity Market (MIBEL) comprising the Spanish and Portuguese electricity systems. The MIBEL market includes in the short-term: the day-ahead market (DAM) and a set of balancing, reserve and adjustment markets (intraday markets); these markets are complemented with the medium- and long-term mechanisms: a derivatives market and different kinds of bilateral contracts. This structure is similar to other European electricity markets and explains why generation companies can no longer optimize their short-term generation planning decisions, i.e. their bidding strategies, without

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considering the relationship between the short-term bid and the medium-term physical products. The MIBEL's directives dictate specific rules describing how these medium-term mechanisms should be included into the DAM bid. This work deals with the most relevant medium-term mechanisms in the MIBEL: the national bilateral contracts (BC) and the future physical contracts (FC). Stochastic programming techniques are applied to maximize the expected value of the utility's profit coming from the day-ahead market, where the significative random variable is the auction clearing price of the day-ahead electricity market. This random variable is modeled through a set of scenarios of the forecasted prices. The set of scenarios is used to feed a two-stage stochastic optimization model that finds the optimal day-ahead bid of a price-taker GenCo (an electrical utility without influence over the market prices) operating in the MIBEL and holding bilateral and physical futures contracts.

The extensive form of the deterministic equivalent of this stochastic programming problem will be a mixed integer quadratic programming problem (MIQP), which is difficult to solve efficiently, particularly for large-scale instances. Several algorithmic approaches can be adopted to overcome this difficulty. In [?] the quadratic objective function of this problem is approximated by a polyhedral outer approximation by means of *perspective cuts* so that we can exploit the efficiency of general-purpose solvers for mixed integer linear problems (MILP). An alternative to the perspective cuts methodology is the Second-Order Cone Program reformulation (SOCP, [?]), but for quadratic problems the perspective cuts reformulation was reported to be more efficient [?]. Finally, the Branch-and-Fix Coordination (BFC) method has been used successfully to solve two-stage stochastic mixed integer linear problems [?] to solve the day-ahead optimal bid problem. In this work we propose an combination between BFC and PC to efficiently solve the optimal day-ahead bid problem.

2 Day-ahead electricity market bid with futures and bilateral contracts model (DAMB-FBC)

In this section the model (DAMB-FBC) is formulated as a two-stage stochastic programming problem that allows a price-taker generation company to optimally decide the unit commitment of its thermal units, the economic dispatch of the bilateral and futures contracts between the thermal units, and the optimal generation bid of the committed units to the MIBEL's day-ahead market. The objective function of the model represents the expected profits of the GenCo's participation in the day-Ahead market. The constraints assure that the MIBEL's rules and the operational restrictions of the units are respected. The main decision variables are the ones that model the start-up and shut-down of the units, the quantity that will be bid at instrumental price and the scheduled energy committed to the bilateral and the futures contracts settlement.

2.1 Parameters

The (DAMB-FBC) model considers a price-taker GenCo owning a set of thermal generation units \mathcal{I} that bid to the $t \in \mathcal{T} = \{1, 2, \dots, 24\}$ hourly auctions of the DAM. The parameters for the i^{th} thermal unit are:

- c_i^b , c_i^l and c_i^q , generation costs with constant, linear and quadratic coefficients (€, €/MWh and €/MWh² respectively).
- \bar{P}_i and \underline{P}_i , upper and lower bounds on the hourly energy generation (MWh).
- c_i^{on} and c_i^{off} , start-up and shut-down costs (€).
- t_i^{on} and t_i^{off} , minimum operation and minimum idle time (h).

A base load physical futures contract $j \in \mathcal{F}$ is defined by:

- \mathcal{U}_j , the set of generation units allowed to cover the FC j .
- L_j^F , the amount of energy (MWh) to be procured each interval of the delivery period by the set \mathcal{U}_j of generation units to cover contract j .
- λ_j^F , the price of contract j (€/MWh).

A base load bilateral contract $k \in \mathcal{B}$ is defined by:

- L_k^B , the amount of energy (MWh) to be procured at each interval of the delivery period by the set of available generation units to cover the BCs.
- λ_k^B , the price of the contract k (€/MWh).

The random variable λ_t^D , the clearing price of the t^{th} hourly auction of the DAM, is represented in the two-stage stochastic model by a set of scenarios $s \in \mathcal{S}$, each one with its associated clearing price for each DAM auction $t \in \mathcal{T}$:

- $\lambda_t^{D,s}$ clearing price for auction t at scenario s (€/MWh).
- P^s probability of scenario s .

2.2 Variables

Those decision variables that doesn't depend on the scenarios are called first stage (or *here-and-now*) variables and in our formulation are, for each $t \in \mathcal{T}$ and $i \in \mathcal{I}$:

- u_{ti} , the unit commitment (binary)
- c_{ti}^u , c_{ti}^d , the start-up/shut-down costs variables.
- q_{ti} , the instrumental price offer bid.
- f_{tij} , the scheduled energy for FC $j \in \mathcal{F}$.
- b_{ti} , the scheduled energy for the pool of BCs .

Decision variables that can adopt different values depending on the scenario are called second stage variables and in our formulation are, for each $t \in \mathcal{T}$, $i \in \mathcal{I}$ and scenario $s \in \mathcal{S}$:

- g_{ti}^s , the total generation.
- p_{ti}^s , the matched energy in the day-ahead market.

2.3 Constraints

Bilateral and futures contracts constraints The coverage of both the physical futures contracts and the bilateral contracts must be guaranteed. The constraints for each futures contract are:

$$\sum_{i \in \mathcal{U}_j} f_{tij} = L_j^F \quad t \in \mathcal{T}, j \in \mathcal{F} \quad (1)$$

$$f_{tij} \geq 0 \quad t \in \mathcal{T}, j \in \mathcal{F}, i \in \mathcal{I} \quad (2)$$

and the bilateral contract constraints are:

$$\sum_{i \in \mathcal{I}} b_{ti} = \sum_{k \in \mathcal{B}} L_k^B \quad t \in \mathcal{T} \quad (3)$$

$$0 \leq b_{ti} \leq \bar{P}_i u_{ti} \quad i \in \mathcal{I}, t \in \mathcal{T} \quad (4)$$

where L_k^B is the energy that has to be settled for contract $k \in \mathcal{B}$

Day-ahead market and total generation constraints As we have introduced, we will use the value of the *matched energy* in our formulation. The matched energy is the accepted energy in the clearing process, that is, the energy generated that will be rewarded at the clearing price. This matched energy is uniquely determined by the sale bid and the clearing price and it will play a central role in the presented model [?].

The MIBEL's rules affecting the day-ahead market establishes the relation between the variables representing the matched energy p_{ti}^s , the energy of the bilateral contracts b_{ti} , the energy of the futures contracts f_{tij} , the instrumental price offer bid q_{ti} , and the commitment binary variables u_{ti} . The energies L_j^F and L_k^B must be integrated in the MIBEL's DAM bid observing the two following rules:

1. If generator i contributes with f_{tij} MWh at period t to the coverage of the FC j , then the energy f_{tij} must be offered to the pool for free (*instrumental price bid*).
2. If generator i contributes with b_{ti} MWh at period t to the coverage of any of the BCs, then the remaining production capacity $\bar{P}_i - b_{ti}$ must be bid to the DAM.

These rules can be included in the model by means of the following set of constraints:

$$p_{ti}^s \geq q_{ti} \quad i \in \mathcal{I}, t \in \mathcal{T}, s \in \mathcal{S} \quad (5)$$

$$p_{ti}^s \leq \bar{P}_i u_{ti} - b_{ti} \quad i \in \mathcal{I}, t \in \mathcal{T}, s \in \mathcal{S} \quad (6)$$

$$q_{ti} \geq \underline{P}_i u_{ti} - b_{ti} \quad i \in \mathcal{I}, t \in \mathcal{T}, s \in \mathcal{S} \quad (7)$$

$$q_{ti} \geq \sum_{j \mid i \in \mathcal{U}_j} f_{tij} \quad i \in \mathcal{I}, t \in \mathcal{T}, s \in \mathcal{S} \quad (8)$$

where:

(??) and (??) ensure respectively that the matched energy p_{ti}^s will be greater than the instrumental price bid q_{ti} and less than the total available energy not allocated to BC.

(??) and (??) guarantee respectively that the instrumental price bid will be greater than the minimum generation output of the unit and greater than the contribution of the unit to the FC coverage.

Please note that (??) together with (??) assures that q_{ti} will be always non-negative. The total generation level of a given unit i , g_{ti}^s , is defined as the addition of the allocated energy to the BC plus the matched energy of the DAM:

$$g_{ti}^s = b_{ti} + p_{ti}^s i \in \mathcal{I}, t \in \mathcal{T}, s \in \mathcal{S} \quad (9)$$

Constraints (??)-(??) assure that g_{ti}^s will be always either zero or $g_{ti}^s \in [\underline{P}_i, \overline{P}_i]$, that is:

$$\underline{P}_i u_{ti} \leq g_{ti}^s \leq \overline{P}_i u_{ti}, i \in \mathcal{I}, t \in \mathcal{T}, s \in \mathcal{S} \quad (10)$$

Unit commitment constraints This section includes the formulation for the unit commitment of the thermal units [?]. The first two sets of constraints model the start-up and shut-down costs and the next ones control minimum operation and idle time for each unit. First, the start-up and shut-down costs are modeled:

$$c_{ti}^u \geq c_i^{on} [u_{ti} - u_{(t-1)i}] \quad i \in \mathcal{I}, t \in \mathcal{T} \setminus \{1\} \quad (11)$$

$$c_{ti}^d \geq c_i^{off} [u_{(t-1)i} - u_{ti}] \quad i \in \mathcal{I}, t \in \mathcal{T} \setminus \{1\} \quad (12)$$

$$c_{ti}^u, c_{ti}^d \geq 0 \quad i \in \mathcal{I}, t \in \mathcal{T} \quad (13)$$

$$u_{ti} \in \{0, 1\} \quad i \in \mathcal{I}, t \in \mathcal{T} \quad (14)$$

The initial state of each thermal unit i can be taken into account through the parameters G_i and H_i that represent, respectively, the number of the initial time periods along which the thermal unit must remain on (G_i) or off (H_i). These conditions are imposed by the following set of constraints:

$$\sum_{j=1}^{G_i} (1 - u_{ji}) = 0 \quad \text{and} \quad \sum_{j=1}^{H_i} u_{ji} = 0, \quad i \in \mathcal{I} \quad (15)$$

Finally, the minimum up and down time, t_i^{on} and t_i^{off} are imposed, up to the periods $|\mathcal{T}| - (t_i^{on} - 1)$ and $|\mathcal{T}| - (t_i^{off} - 1)$, through the following set of constraints:

$$\sum_{n=t}^{t+t_i^{on}-1} u_{in} \geq t_i^{on} [u_{ti} - u_{(t-1)i}] \quad t = G_i + 1, \dots, |\mathcal{T}| - t_i^{on} + 1, i \in \mathcal{I} \quad (16)$$

$$\sum_{n=t}^{t+t_i^{off}-1} (1 - u_{ni}) \geq t_i^{off} [u_{(t-1)i} - u_{ti}] \quad t = H_i + 1, \dots, |\mathcal{T}| - t_i^{off} + 1, i \in \mathcal{I} \quad (17)$$

and for the last $t_i^{on} - 1$ and $t_i^{off} - 1$ time periods:

$$\sum_{n=t}^{|\mathcal{T}|} (u_{ni} - [u_{ti} - u_{(t-1)i}]) \geq 0 \quad t = |\mathcal{T}| - t_i^{on} + 2, \dots, |\mathcal{T}|, i \in \mathcal{I} \quad (18)$$

$$\sum_{n=t}^{|\mathcal{T}|} (1 - u_{ni} - [u_{(t-1)i} - u_{ti}]) \geq 0 \quad t = |\mathcal{T}| - t_i^{off} + 2, \dots, |\mathcal{T}|, i \in \mathcal{I} \quad (19)$$

2.4 Objective function

The quadratic function that gives the long-run expected profits of the GenCo after the participation in the DAM is:

$$\min E_{\lambda^D} [C(u, c^u, c^d, g, p; \lambda^D)] = \sum_{t \in \mathcal{T}} \sum_{i \in \mathcal{I}} (c_{ti}^u + c_{ti}^d + c_{ti}^b u_{ti} + \quad (20)$$

$$+ \sum_{s \in \mathcal{S}} P^s [(c_{ti}^l g_{ti}^s + c_{ti}^q (g_{ti}^s)^2) - \lambda_t^{D,s} p_{ti}^s]), \quad (21)$$

where the right hand side of (??) is the on/off fixed cost of the unit commitment of the thermal units, deterministic and independent of the realization of the random variable $\lambda_t^{D,s}$ and (??) represents the expected value of the benefits from the DAM. The term between parenthesis corresponds to the expression of the quadratic generation costs associated to the total generation of the unit g_{ti}^s while the last term, $\lambda_t^{D,s} p_{ti}^s$ computes the incomes from the DAM due to a value p_{ti}^s of the matched energy.

Please note that the constant incomes from the BC and FC, i.e. $\sum_{k \in \mathcal{B}} \lambda_k^{BC} L_k^{BC}$

and $\sum_{t \in \mathcal{T}, j \in \mathcal{J}} (\lambda_j^{FC} - \bar{\lambda}_t^D) L_j^{FC}$, have been dropped from the objective function.

2.5 Model (DAMB-FBC)

The model defined so far can be represented as:

$$(DAMB-FBC) \left\{ \begin{array}{ll} \min E_{\lambda^D} [C(u, c^u, c^d, g, p; \lambda^D)] & \\ \text{s.t.} & \\ \text{Eq. (??) - (??)} & \text{BC and FC constraints} \\ \text{Eq. (??) - (??)} & \text{DAM and total gen. constraints} \\ \text{Eq. (??) - (??)} & \text{Unit commitment constraints} \end{array} \right.$$

Model (DAMB-FBC) is the optimization problem associated with the two-stage stochastic programming problem with a set \mathcal{S} of scenarios for the spot price λ_t^D , where $t \in \mathcal{T}$. This optimization problem is a convex MIQP with a well defined global optimal solution.

3 QBFC method

Model (DAMB-FBC) can be rewritten as the so-called Deterministic Equivalent Model (DEM)

$$\begin{aligned}
& \text{minimize } c^t \delta + \sum_{s \in \mathcal{S}} P^s q^s(x, y^s) \\
& \text{subject to : } l_a \leq A \begin{bmatrix} \delta \\ x \end{bmatrix} \leq u_a, \\
& \quad l_t^s \leq T^s \begin{bmatrix} \delta \\ x \\ y^s \end{bmatrix} \leq u_t^s, \quad s \in \mathcal{S}, \\
& \quad x \geq 0, \underline{y} \leq y^s \leq \bar{y}, \quad s \in \mathcal{S}, \\
& \quad \delta \in \{0, 1\}^{n_\delta},
\end{aligned}$$

where $\delta = u$, $x = (c^u, c^d)$, $y = (g, p)$, $q(x, y) = b_x^t x + b_y^t y + y^t Q_{yy} y$, and Q_{yy} being a diagonal matrix.

As is showed by [?] the compact representation (DEM) can be written as a *splitting variable* representation; i.e., δ and x are respectively replaced by δ^s and x^s , for $s \in \mathcal{S}$. So, we have

$$\begin{aligned}
(\text{MIQP}) \quad & \text{minimize } \sum_{s \in \mathcal{S}} P^s (c^t \delta^s + q^s(x^s, y^s)) \\
& \text{subject to : } l_a \leq A \begin{bmatrix} \delta^s \\ x^s \end{bmatrix} \leq u_a, \quad s \in \mathcal{S}, \\
& \quad l_t^s \leq T^s \begin{bmatrix} \delta^s \\ x^s \\ y^s \end{bmatrix} \leq u_t^s, \quad s \in \mathcal{S}, \\
& \quad x^s \geq 0, \underline{y} \leq y^s \leq \bar{y}, \quad \delta^s \in \{0, 1\}^{n_\delta}, \quad s \in \mathcal{S}, \\
& (\text{NAC}_\delta) \quad \delta^s - \delta^{s'} = 0, \quad \forall s, s' \in \mathcal{S} : s \neq s', \\
& (\text{NAC}_x) \quad x^s - x^{s'} = 0, \quad \forall s, s' \in \mathcal{S} : s \neq s',
\end{aligned}$$

where NAC_δ and NAC_x are the *nonanticipativity constraints*.

In this method (DEM) is solved by using a Branch-and-Fix-Coordination scheme (BFC) for each scenario $s \in \mathcal{S}$ to fulfill the integrality condition (IC) on the variables δ , so that the NAC_δ are also satisfied when selecting branching nodes and branching variables by the Twin-Node-Families concept (TNF), which was introduced by [?].

A similar approach to that suggested in [?] is used in this work to coordinate the selection of the branching node and branching variable for each scenario-related BF tree, such that the NAC_δ are satisfied when fixing δ^s , for all $s \in \mathcal{S}$, either to 1 or to 0. A *TNF integer set* is a set of integer BF nodes (i.e. they verify IC), one per BF tree, in which the NAC_δ are verified. More details about this methodology can be found in [?].

When the number of scenarios is very high, in order to gain computational efficiency we can take scenario clusters; i.e., instead a submodel for each scenario $s \in \mathcal{S}$ we can use a submodel (MIQP^p) for each scenario cluster $\mathcal{S}^p \subset \mathcal{S}$ with $p = 1, \dots, \widehat{p}$, where $\mathcal{S}^p \cap \mathcal{S}^{p'} = \emptyset$, for all $p \neq p'$, and $\bigcup_{p=1}^{\widehat{p}} \mathcal{S}^p = \mathcal{S}$,

$$\text{(MIQP}^p\text{)} \quad \text{minimize} \quad \sum_{s \in \mathcal{S}^p} P^s (c^t \delta^p + q^s(x^p, y^s)), \quad (22)$$

$$\text{subject to: } l_a \leq A \begin{bmatrix} \delta^p \\ x^p \end{bmatrix} \leq u_a, \quad (23)$$

$$l_t^s \leq T^s \begin{bmatrix} \delta^p \\ x^p \\ y^s \end{bmatrix} \leq u_t^s, s \in \mathcal{S}^p, \quad (24)$$

$$x^p \geq 0, \underline{y} \leq y^s \leq \bar{y}, s \in \mathcal{S}^p, \quad \delta^p \in \{0, 1\}^{n_\delta}, \quad (25)$$

These submodels are linked by the NACs $\delta^p - \delta^{p'} = 0$ and $x^p - x^{p'} = 0$, for all $p, p' \in \{1, \dots, \widehat{p}\}$ such that $p \neq p'$.

In order to gain computational efficiency we propose to use perspective cuts (PC) [?,?] to solve the quadratic subproblems in each node of the TNF. Then MIQP^p becomes:

$$\begin{aligned} \min \quad & \sum_{s \in \mathcal{S}^p} P^s \left\{ \left(b_x^t x + x^t Q_{xx} x \right) + \left(\sum_{i=1}^n v_i^s \right) \right\} \\ \text{s.t.:} \quad & v_i^s \geq (2q_{ii}^s \underline{y}_i + b_i^s) y_i^s + (c_i - q_{ii}^s \underline{y}_i^2) \delta_i^s, \quad i \in \{1, \dots, n\}, s \in \mathcal{S}^p, \\ & v_i^s \geq (2q_{ii}^s \bar{y}_i + b_i^s) y_i^s + (c_i - q_{ii}^s \bar{y}_i^2) \delta_i^s, \quad i \in \{1, \dots, n\}, s \in \mathcal{S}^p, \\ & \text{Eq. (??) - (??)}. \end{aligned}$$

These methods have been implemented in C++ with the help of Cplex 12.1 to solve only the quadratic subproblems. In this work two algorithmic alternatives have been considered:

- ▷ QBFC: coordination of δ in the TNF of the BF trees for clusters $p \in \{1, \dots, \widehat{p}\}$ without using PCs.
- ▷ QBFC-PC: coordination of δ in the TNF of the BF trees for clusters $p \in \{1, \dots, \widehat{p}\}$ using PCs.

For our instances the number of scenarios in each cluster is the same, $|\mathcal{S}^p| = |\mathcal{S}|/\widehat{p}$. Each cluster contains $|\mathcal{S}^p|$ consecutive scenarios, starting from the first one and following in natural order.

4 Numerical Tests

These instances are based on the liberalized electricity market model suggested in [?]. In these problems Q_{xx} is the zero matrix, as a result, when we use perspective cuts the subproblem to solve in each node is linear. The tests have been

performed on HP with Intel(R) Core(TM)2 Quad CPU Q8300 2.50GHz 4 CPU under SUSE Linux Enterprise Desktop 11 (x86_64).

In Table ?? $|\mathcal{S}|$ means the number of scenarios, $|\mathcal{T}|$ the number of periods, “# var” the number of continuous variables, “# var_{PCF}” the number of continuous variables for the PC formulation, “# bin” the number of binary variables, and “# constr” the number of constraints for (DEM).

Table 1. Test problems

Prob.	$ \mathcal{S} $	$ \mathcal{T} $	# var	# var _{PCF}	# bin	# constr
P01	10	12	1296	1776	48	1788
P02	20	12	2256	3216	48	3228
P03	30	12	3216	4656	48	4668
P04	40	12	4176	6096	48	6108
P05	50	12	5136	7536	48	7548
P11	10	24	2592	3552	96	3600
P12	20	24	4512	6432	96	6480
P13	30	24	6432	9312	96	9360
P14	40	24	8352	12192	96	12240
P15	50	24	10272	15072	96	15120

For every problem $|\mathcal{F}| = |\mathcal{B}| = 2$ and $|\mathcal{I}| = 4$. If we use the PC formulation, the problem increases the number of variables in $m = |\mathcal{T}| \cdot |\mathcal{I}| \cdot |\mathcal{S}|$ and the number of constraints in $2 \cdot m$.

Table 2. Computational results: CPU-times

Prob.	\hat{p}	QBFC	QBFC-PC	ratio	# PC
P01	2	10.1	3.4	0.34	280
P02	4	18.7	8.7	0.47	825
P03	5	2153.0	39.8	0.02	1685
P04	5	50.0	45.1	0.90	1491
P05	5	113.7	19.5	0.17	1276
P11	2	86.8	27.4	0.32	513
P12	4	469.7	50.3	0.11	1821
P13	5	687.3	176.6	0.26	3454
P14	5	1198.0	276.7	0.23	4239
P15	5	1190.9	246.3	0.21	2592

In Table ?? below the headings QBFC are the times in CPU-seconds used for solving problems with the number of scenario cluster given below the heading \hat{p} and by solving the quadratic subproblem QP^p for each node using Cplex. Column QBFC-PC gives us the CPU-seconds and indicates that the quadratic subproblems QP^p have been solved by using perspective cuts, which means that instead of solving a quadratic problem QP^p in each node of a TNF for $p \in$

$\{1, 2, \dots, \widehat{p}\}$, a linear problem is solved. Also, “ratio” = $\frac{\text{QBFC-PC}}{\text{QBFC}}$ gives us the ratio of CPU-times. Note that the running time with PC is a 30% of the running time without PC (average). The last column, “# PC”, means the number of perspective cuts generated in each test.

5 Conclusions

We have presented an Optimal Bidding Model for a price-taker generation company operating both in the MIBEL Derivatives and Day-Ahead Electricity Market (DAMB-FBC). The model developed finds the optimal bid for the spot market, the optimal allocation of the physical futures and bilateral contracts among the thermal units and the unit commitment following in detail the MIBEL rules. The (DAMB-FBC) has been solved both with the standard BFC method and with a PC variation which reduces the running time to a 30% on the average.

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