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Title	The Tate-Lichtenbaum pairing on a hyperelliptic curve via hyperelliptic nets
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Citation	(2011), 2011: 151-151
Issue Date	2011
URL	http://hdl.handle.net/2433/214946
Right	
Type	Departmental Bulletin Paper
Textversion	publisher

Kyoto University

The Tate-Lichtenbaum pairing on a hyperelliptic curve via hyperelliptic nets
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1 Introduction

Let C be a non-singular projective curve defined over a finite field $\mathbb{F}_q$. Let m be a positive integer with $m \mid (q-1)$. The Tate-Lichtenbaum pairing is a bilinear map

$$\tau_m \colon \operatorname{Pic}^0(C)[m] \times \operatorname{Pic}^0(C)/m\operatorname{Pic}^0(C) \to \mathbb{F}_q^\times/(\mathbb{F}_q^\times)^m.$$

For cryptographic applications, it is important to find an efficient algorithm to compute the Tate-Lichtenbaum pairing.

The Tate-Lichtenbaum pairing is usually computed by Miller's algorithm. Recently, Stange [3] gave a new algorithm to compute the Tate(-Lichtenbaum) pairing on an elliptic curve. This algorithm is based on elliptic nets, which are also defined by Stange as a generalization of elliptic divisibility sequences.

In this poster, we define hyperelliptic nets as a generalization of elliptic nets to hyperelliptic curves. We also give an expression for the Tate-Lichtenbaum pairing on a hyperelliptic curve in terms of hyperelliptic nets. By using this expression, we obtain an algorithm to compute the Tate-Lichtenbaum pairing on a hyperelliptic curve of genus 2.

2 The hyperelliptic sigma function

Let C be a non-singular projective curve of genus g over $\mathbb{C}$ defined by

$$y^2 = x^{2g+1} + \lambda_{2g}x^{2g} + \cdots + \lambda_1 x + \lambda_0.$$

We use the following notation:
- ∞: the point at infinity of C,
- J: the Jacobian variety of C,
- $\kappa \colon \mathbb{C}^g \to \mathbb{C}^g/\Lambda \cong J(\mathbb{C})$, where we fix a uniformization $\mathbb{C}^g/\Lambda \cong J(\mathbb{C})$,
- $\lambda \colon C \to J$: an embedding with $\lambda(\infty) = O$,
- $\Theta = \lambda(C) + \cdots + \lambda(C)$ ($g-1$ times): the theta divisor.

We write $e(z) = \exp(2\pi\sqrt{-1}\,z)$. We define the theta function with characteristics by

$$\vartheta\begin{bmatrix} a \\ b \end{bmatrix}(z,\tau) = \sum_{n \in \mathbb{Z}^g} e\left(\frac{1}{2}\,{}^t(n+a)\tau(n+a) + {}^t(n+a)(z+b) \right),$$

where $z \in \mathbb{C}^g$, $\tau \in M_g(\mathbb{C})$ is symmetric, $\operatorname{Im}(\tau)$ is positive definite, and $a, b \in \mathbb{R}^g$.

Definition 1 (cf. [2]). We define the *hyperelliptic sigma function* on $\mathbb{C}^g$ by

$$\sigma(u) = c \exp\left(\frac{1}{2}\,{}^t u \eta' \omega'^{-1} u \right) \vartheta[\delta](\omega'^{-1}u, \omega'^{-1}\omega''),$$

where $\omega', \omega'' \in M_g(\mathbb{C})$ are period matrices satisfying $\Lambda = \omega'\mathbb{Z}^g + \omega''\mathbb{Z}^g$ and $c \in \mathbb{C}$, $\eta' \in M_g(\mathbb{C})$, and $\delta \in ((1/2)\mathbb{Z})^{2g}$ are certain constants.

If $g = 1$, then $\sigma(u)$ coincides with the Weierstrass sigma function. The sigma function $\sigma(u)$ has a zero of order 1 along $\kappa^{-1}(\Theta)$.

3 Hyperelliptic nets

Let n be a positive integer.

We first define hyperelliptic nets over $\mathbb{C}$. The following is a straightforward generalization of Stange's definition.

Definition 2. Let $P_1, \ldots, P_n \in J(\mathbb{C}) \setminus \Theta$ with $P_i + P_j \notin \Theta$ for all $1 \le i < j \le n$. Let $u_1, \ldots, u_n \in \mathbb{C}^g$ be points with $\kappa(u_i) = P_i$. We define a map $W_{C,P_1,\ldots,P_n} \colon \mathbb{Z}^n \to \mathbb{C}$ by

$$W_{C,P_1,\ldots,P_n}(v_1,\ldots,v_n) = \frac{\sigma(v_1 u_1 + \cdots + v_n u_n)}{\prod_{i=1}^n \sigma(u_i)^{2v_i^2 - \sum_{j=1}^n v_i v_j} \prod_{1 \le i < j \le n} \sigma(u_i + u_j)^{v_i v_j}}.$$

We call $W_{C,P_1,\ldots,P_n}$ the *hyperelliptic net associated to $C, P_1, \ldots, P_n$*.

It is easily verified that $W_{C,P_1,\ldots,P_n}$ does not depend on the choice of $u_1, \ldots, u_n$.

Lemma 3. We assume that C is defined over a subfield $K \subset \mathbb{C}$ and $P_1, \ldots, P_n \in J(K)$. Then, for any $v \in \mathbb{Z}^n$,

$$W_{C,P_1,\ldots,P_n}(v) \in K.$$

By Lemma 3, we can define hyperelliptic nets over any field of characteristic 0. Furthermore, we can also define hyperelliptic nets over any field of positive characteristic by reduction.

4 Recurrence relations

Theorem 4. Let $W_{C,P_1,\ldots,P_n}$ be a hyperelliptic net as before. Let $m > 2^g$ be an integer. Let $v^{(1)}, \ldots, v^{(m)} \in ((1/2)\mathbb{Z})^n$ with $v^{(i)} + v^{(j)}, v^{(i)} - v^{(j)} \in \mathbb{Z}^n$ for all $1 \le i, j \le m$. We define a matrix A by

$$A = \left(W_{C,P_1,\ldots,P_n}(v^{(i)} + v^{(j)}) W_{C,P_1,\ldots,P_n}(v^{(i)} - v^{(j)}) \right)_{1 \le i,j \le m}.$$

Then we have $\det A = 0$. In particular, if $g \equiv 1, 2 \pmod 4$ and m is even, then we have $\operatorname{pf} A = 0$, where $\operatorname{pf} A$ is the Pfaffian of A.

5 The Tate-Lichtenbaum pairing

Definition 5. The *Tate-Lichtenbaum pairing* is a map

$$\tau_m \colon \operatorname{Pic}^0(C)[m] \times \operatorname{Pic}^0(C)/m\operatorname{Pic}^0(C) \to \mathbb{F}_q^\times/(\mathbb{F}_q^\times)^m;$$

$$(\overline{D}, \overline{E}) \mapsto f_D(E) = \prod_{i=1}^r f_D(P_i)^{n_i},$$

where D and E are divisors on C representing $\overline{D}$ and $\overline{E}$ respectively such that D and E have no points in common, f_D is a rational function on C with $\operatorname{div}(f_D) = mD$, and $E = \sum_{i=1}^r n_i P_i$.

It is known that τ_m is bilinear and non-degenerate (cf. [1]).

The Tate-Lichtenbaum pairing is described in terms of hyperelliptic nets.

Theorem 6. Let P and Q be points associated to $\overline{D}$ and $\overline{E}$ respectively. Then we have

$$\tau_m(\overline{D}, \overline{E}) = \frac{W_{C,P,Q}(m+1,1) W_{C,P,Q}(1,0)}{W_{C,P,Q}(m+1,0) W_{C,P,Q}(1,1)} \bmod (\mathbb{F}_q^\times)^m.$$

Theorem 6 is a generalization of Stange's result for elliptic curves.

By Theorems 4 and 6, we can compute $\tau_m(\overline{D}, \overline{E})$ for a curve of genus 2.

Corollary 7. Assume $g = 2$. If (P, Q) is not on a certain divisor on $J \times J$, then $\tau_m(\overline{D}, \overline{E})$ can be computed with $O(\log m)$ operations in $\mathbb{F}_q$.

Note that the above complexity is the same as that of Miller's algorithm.

6 Example

Let $q = 47$. We consider the following curve and divisors:
- $C/\mathbb{F}_{47} \colon y^2 = x^5 + x + 41$,
- $D = (\alpha_1, 31\alpha_1 + 3) + (\alpha_2, 31\alpha_2 + 3) - 2\infty$,
- $E = (\beta_1, 22\beta_1 + 14) + (\beta_2, 22\beta_2 + 14) - 2\infty$,

where α_1 and α_2 are the roots of the polynomial $x^2 + 6x + 16$ over $\mathbb{F}_{47}$, and β_1 and β_2 are the roots of the polynomial $x^2 + 29x + 24$ over $\mathbb{F}_{47}$.

Let P and Q be the points on J corresponding to $\overline{D}$ and $\overline{E}$ respectively. The following figure shows a part of the net $W_{C,P,Q}$.

	7	12	2	36	19	33	39	18
	14	38	14	10	23	21	36	9
	13	31	32	18	8	2	2	16
	14	15	43	19	44	5	22	42
	25	6	33	11	10	36	21	16
	25	8	2	13	16	32	14	5
↑	1	1	23	4	29	40	43	7
Q	0	1	37	18	36	2	7	45

$P \to$

Let $m = 23$. Then m is a divisor of $q - 1 = 46$ and $\overline{D} \in \operatorname{Pic}^0(C)[m]$. We have

$$W_{C,P,Q}(m+1,1) = 43, \quad W_{C,P,Q}(m+1,0) = 8,$$
$$W_{C,P,Q}(1,0) = W_{C,P,Q}(1,1) = 1.$$

Therefore, by Theorem 6,

$$\tau_m(\overline{D}, \overline{E}) = \frac{43}{8} \bmod (\mathbb{F}_{47}^\times)^{23} = 23 \bmod (\mathbb{F}_{47}^\times)^{23}.$$

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