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## Long Cycles through Specified Vertices in a Graph

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### ABSTRACT

In this paper, we consider the length of the longest cycle through specified vertices. We show the following two results. (1) Let  $G$  be a  $k$ -connected graph of order at least  $2k$  and circumference  $l$ . Suppose  $m < k$ . Then for any  $m$  vertices of  $G$ ,  $G$  has a cycle which contains all of them and has length at least  $\frac{k-m}{k}l + 2m$ . (2) Let  $G$  be a 3-connected planar graph with circumference  $l$ . Then for any three vertices of  $G$ , there exists a cycle which contains all of them and has length at least  $\frac{1}{4}l + 3$ .

Here, we consider finite simple graphs. Let  $G$  be a graph. By Dirac's theorem[3]  $G$  has a cycle through specified  $k$  vertices. In [2] Dirac also showed that a 2-connected graph of order  $n$  and minimum degree at least  $d$  has a cycle of length at least  $\min\{n, 2d\}$ . Locke[4] and Voss[7] generalized his result by showing that under the same conditions the graph has a cycle of length at least  $\min\{n, 2d\}$  which contains specified two vertices.

These results lead us to the following question: Does a  $k$ -connected graph have a long cycle through specified  $m$  vertices ( $m \leq k$ )? In this paper we investigate this question.

For basic graph-theoretic terminology, we refer the reader to [1]. Let  $G$  be a graph. The *circumference* of  $G$ , denoted by  $\text{cir}(G)$ , is the length of the longest cycle of  $G$ . We denote by  $w(G)$  the number of components of  $G$ . For  $k \geq 0$  and  $S \subset V(G)$ , we call  $S$  a  $k$ -cutset if  $w(G - S) \geq 2$  and  $|S| = k$ . We often identify a subgraph  $H$  of  $G$  with its vertex set  $V(H)$ . Especially, when  $x$  is a vertex of  $H$ , we write  $x \in H$  instead of  $x \in V(H)$ . Furthermore, we write  $|H|$  instead of  $|V(H)|$ . When we consider a cycle, we always give it an orientation. Let  $C^+$  be the orientation of a cycle  $C$  and  $C^-$  be its reverse orientation. Let  $C^+ = x_0, x_1, \dots, x_{n-1}, x_n$  be a cycle. For  $x_i, x_j \in C$ , we define a subpaths  $C^+[x_i, x_j]$  and  $C^-[x_i, x_j]$  of  $C$  by

$$C^+[x_i, x_j] = x_i, x_{i+1}, \dots, x_{j-1}, x_j,$$

and

$$C^-[x_i, x_j] = x_i, x_{i-1}, \dots, x_{j+1}, x_j.$$

We also define  $C^+(x_i, x_j)$  and  $C^-(x_i, x_j)$  by

$$C^+(x_i, x_j) = C^+[x_i, x_j] - \{x_i, x_j\},$$

and

$$C^-(x_i, x_j) = C^-[x_i, x_j] - \{x_i, x_j\}.$$

Furthermore,  $C^+[x_i, x_j] = C^+[x_i, x_j] - \{x_j\}$ . Subpaths  $C^-[x_i, x_j]$ ,  $C^+(x_i, x_j)$ ,  $C^-(x_i, x_j)$  are defined similarly. Let  $x_1, x_2, \dots, x_s$  be a path. We denote by  $\text{end}(P)$  the set of endvertices of  $P$ ;  $\text{end}(P) = \{x_1, x_s\}$ . Let  $P = x_1, x_2, \dots, x_s$  and  $Q = y_1, y_2, \dots, y_t$  be paths such that  $x_s = y_1$ . We denote by  $P \cdot Q$  the walk  $x_1, x_2, \dots, x_s = y_1, y_2, \dots, y_t$ .

Let  $z \in V(G)$  and  $S \subset V(G) - \{z\}$ . A subgraph  $F$  of  $G$  is called a  $(z, S)$ -fan if  $F$  has the following decomposition  $F = \cup_{i=1}^k P_i$ , where

- (1) each  $P_i$  is a path between  $z$  and  $a_i \in S$ , and
- (2)  $P_i \cap S = \{a_i\}$ , and  $P_i \cap P_j = \{z\}$  if  $i \neq j$ .

We call  $k$  the size of the fan  $F$ . The vertices  $a_1, \dots, a_k$  are called endvertices of  $F$  and the set of its endvertices is denoted by  $\text{end}(F)$ . Since  $F$  is a tree, for any two vertices  $x, y \in F$  the path in  $F$  which joins  $x$  and  $y$  is unique. We denote this path by  $F[x, y]$ . We define  $F[x, y]$  by  $F[x, y] = F[x, y] - \{y\}$ . Paths  $F(x, y)$  and  $F(x, y)$  are defined similarly.

The following theorem is well-known, called the generalized Menger's theorem.

**THEOREM A** ([1, Theorem 6.7]). *Let  $G$  be a  $k$ -connected graph,  $z \in V(G)$ , and  $S \subset V(G) - \{z\}$ . Then  $G$  has a  $(z, S)$ -fan of size  $\min\{|S|, k\}$ . ■*

The following theorem was proved by Perfect[5].

**THEOREM B** (Perfect[5]). *Let  $G$  be a graph,  $z \in V(G)$ , and  $S \subset V(G) - \{z\}$ . Suppose  $G$  has two  $(z, S)$ -fans  $F_1$  and  $F_2$  of size  $k_1$  and  $k_2$ , respectively. If  $k_1 \leq k_2$ , then  $G$  has a  $(z, S)$ -fan  $F'$  of size  $k_2$  such that  $\text{end}(F_1) \subset \text{end}(F')$ . ■*

We use these two theorems in the proofs our results.

First, we show that the existence of long cycles through specified  $m$  vertices in a  $k$ -connected graph is assured if  $m < k$ . Note that a  $k$ -connected graph is hamiltonian if its order is at most  $2k$ , by Dirac's theorem.

**THEOREM 1.** Let  $k \geq 2$ ,  $0 \leq m \leq k$  and  $G$  be a  $k$ -connected graph of order at least  $2k$ . For any  $m$  vertices  $x_1, \dots, x_m$  of  $G$ , there exists a cycle such that

- (1)  $x_1, \dots, x_m \in V(C)$ , and
- (2)  $|C| \geq \frac{k-m}{k} \text{cir}(G) + 2m$ .

Recently, Seymour and Truemper sent me a proof which is simpler than the original one. We show their proof.

**Proof** (due to Seymour and Truemper). The proof is by induction on  $m$ . For  $m = 1$ , let  $x \in V(G)$ , and let  $C$  be a longest cycle in  $G$ . Since  $|C| \geq 2k$ ,

$$\frac{k-1}{k} \text{cir}(G) + 2 = |C| - \frac{|C|}{k} + 2 \leq |C|.$$

So we may assume  $x \notin V(C)$ . Now  $G$  has an  $(x, C)$ -fan of size  $k$ . The endvertices of  $F$  divide  $C$  into  $k$  paths, and any shortest one  $P$  of these paths, say  $P = C^+[u, v]$  has length at most  $\frac{1}{k} \text{cir}(G)$ . So  $C^+[v, u] \cdot F[u, v]$  is a cycle which contains  $x$  and has length at least

$$|C| - \frac{\text{cir}(G)}{k} + 2 = \frac{k-1}{k} \text{cir}(G) + 2$$

as desired.

Suppose  $m > 1$ , and let  $C$  be a longest cycle containing at least  $m-1$  members of  $S$ . By the induction hypothesis,

$$\begin{aligned} |C| &\geq \frac{k-m+1}{k} \text{cir}(G) + 2(m-1) \\ &= \frac{k-m}{k} \text{cir}(G) + 2m + \frac{\text{cir}(G)}{k} - 2 \\ &\geq \frac{k-m}{k} \text{cir}(G) + 2m. \end{aligned} \quad (*)$$

So we may assume that exactly one member  $x$  of  $S$  does not lie on  $C$ . Since  $\text{cir}(G) \geq 2k$ ,  $|C| \geq 2k$ . So  $G$  has an  $(x, C)$ -fan of size  $k$ . The endvertices of  $F$  divide  $C$  into  $k$  paths. We call such a path *bad* if it contains some member of  $S$  internally, and we call it *good* if it is not bad. Let  $b$  represent the number of bad paths, and let  $L$  be the sum of lengths of the bad paths. Then some good path  $P = C^+[u, v]$  has length at most

$$\frac{|C| - L}{k - b}$$

(, where  $|C| \geq 2k$  and  $k \geq m-1$ ). Keeping  $|C|$  and  $k$  fixed, and under the conditions  $L \geq 2b$  and  $b \leq m-1$ , this is maximized when  $L = 2b$  and  $b = m-1$ . Hence,

$$|P| \leq \frac{|C| - 2(m-1)}{k - m + 1}.$$

A cycle  $C^+[v, u] \cdot F[u, v]$  contains  $S$ , and from (\*) it has length at least

$$|C| - \frac{|C| - 2(m-1)}{k-m+1} + 2 \geq \frac{k-m}{k} \text{cir}(G) + 2m$$

as desired. ■

Theorem 1 is sharp. Let,  $k \geq 2$ ,  $s \geq 1$ , and  $0 \leq m \leq k$ . Let  $H_0, H_1, \dots, H_k$  and  $H'_0$  be graphs such that  $H_1 \simeq \dots \simeq H_k \simeq K_s$ ,  $H_0 \simeq \overline{K_m}$  and  $H'_0 \simeq \overline{K_k}$ . Suppose vertex sets  $V(H_0), \dots, V(H_k)$  and  $V(H'_0)$  are disjoint. Define  $G(k, m, s)$  by  $G(k, m, s) = (H_1 \cup \dots \cup H_k \cup H_0) + H'_0$ . Then  $G(k, m, s)$  is  $k$ -connected,  $|G(k, m, s)| = ks + k + m \geq 2k$ , and  $\text{cir}(G(k, m, s)) = ks + k$ . On the other hand, the length of the longest cycle through  $V(H_0)$  is  $(k-m)s + k + m$ . The above example shows that large circumference does not assure the existence of long cycles through specified  $k$  vertices in  $k$ -connected graphs.

Next, we confine ourselves to planar graphs. Even if we consider only planar graphs, the length of the longest cycle through specified two vertices in a 2-connected graph is independent of its circumference. Let  $C = x_0, x_1, \dots, x_m = x_0$  be a cycle of length  $m$  ( $m \geq 4$ ). Add a new vertex  $y$  and join  $yx_1$  and  $yx_{m-1}$ . Then this graph has circumference  $m$ , but the unique cycle through  $y$  and  $x_0$  has length four. On the other hand, by Tutte's theorem[6] 4-connected planar graphs are hamiltonian, and hence the length of the longest cycle through four specified vertices in a 4-connected planar graph is equal to its circumference. On a planar graph of connectivity three, we show the following theorem.

**THEOREM 2.** *Let  $G$  be a 3-connected planar graph. Then any three vertices of  $G$  lie on a cycle of length at least  $\frac{1}{4}\text{cir}(G) + 3$ .*

The proof of Theorem 2 is given by the following two lemmas.

**LEMMA 1.** *Let  $G$  be a 3-connected planar graph. Then for any two vertices  $x, y$ , there exists a cycle  $C$  such that*

- (1)  $x, y \in V(C)$ .
- (2)  $|C| \geq \frac{1}{2}\text{cir}(G) + 2$ .

**LEMMA 2.** *Let  $G$  be a 3-connected planar graph,  $x, y, z \in V(G)$  and  $C$  be a cycle of  $G$  such that  $x, y \in V(C)$ . Then there exists a cycle  $C'$  such that*

- (1)  $x, y, z \in V(C')$ .
- (2)  $|C'| \geq \frac{1}{2}|C| + 2$ .

**Proof of Lemma 1.** If  $G$  is hamiltonian, then the lemma clearly holds. So we may assume that  $G$  is not hamiltonian, which implies  $|G| \geq 7$  and  $\text{cir}(G) \geq 6$ . Let  $C$  be a longest cycle of  $G$ . We consider three cases.

*Case 1.*  $\{x, y\} \subset V(C)$ .

This case is trivial.

*Case 2.*  $|\{x, y\} \cap V(C)| = 1$ .

We may assume that  $x \in V(C)$  and  $y \notin V(C)$ . Consider a  $(y, C)$ -fan  $F$  of size three. Let  $\text{end}(F) = \{y_1, y_2, y_3\}$ . If  $x \in \{y_1, y_2, y_3\}$ , say  $x = y_1$ , then we have two cycles  $C^+[x, y_2] \cdot F[y_2, x]$  and  $C^- [x, y_2] \cdot F[y_2, x]$ , one of which has length at least  $\frac{1}{2}|C| + 2 = \frac{1}{2}\text{cir}(G) + 2$  and contains both  $x$  and  $y$ . Next, assume  $x \notin \{y_1, y_2, y_3\}$ . We may assume  $x \in C^+(y_3, y_1)$ . Then one of the two cycles  $C^+[y_3, y_2] \cdot F[y_2, y_3]$  and  $C^- [y_1, y_2] \cdot F[y_2, y_1]$  has the desired properties.

*Case 3.*  $\{x, y\} \cap V(C) = \emptyset$ .

First, we show the following claims.

*Claim 1.* Suppose there exists a path  $P$  in  $G$  such that

- (1)  $P$  joins two distinct vertices of  $C$  and  $P$  intersects  $C$  only at its endvertices.
- (2)  $x, y \in V(P)$ .

Then the Lemma follows.

*Proof.* Let  $a$  and  $b$  be endvertices of  $P$ . Then one of the two cycles  $P[a, b] \cdot C^+[b, a]$  and  $P[a, b] \cdot C^- [b, a]$  satisfies the desired properties.

*Claim 2.* Suppose there exist two paths  $P$  and  $Q$  such that

- (1)  $V(P) \cap V(Q) = \emptyset$ .
- (2) Both  $P$  and  $Q$  join two vertices of  $C$ .
- (3)  $V(P) \cap V(C) = \text{end}(P)$  and  $V(Q) \cap V(C) = \text{end}(Q)$ .
- (4) Vertices of  $\text{end}(P)$  and vertices of  $\text{end}(Q)$  appear alternately around  $C^+$ .
- (5)  $x \in V(P)$  and  $y \in V(Q)$ .

Then the lemma follows.

*Proof.* Let  $\text{end}(P) = \{x_1, x_2\}$  and  $\text{end}(Q) = \{y_1, y_2\}$ . We may assume  $x_1, y_1, x_2$  and  $y_2$  appear in this order around  $C^+$ . Then one of the two cycles

$$C^+[x_1, y_1] \cdot Q[y_1, y_2] \cdot C^- [y_2, x_2] \cdot P[x_2, x_1]$$

and

$$C^- [x_1, y_2] \cdot Q[y_2, y_1] \cdot C^+[y_1, x_2] \cdot P[x_2, x_1]$$

has the desired properties.

Let  $\text{end}(F_1) = \{x_1, x_2, x_3\}$ . We may assume that  $x_1, x_2, x_3$  appear in this order around  $C^+$ . If  $y \in V(F_1)$ , then the theorem follows by Claim 1. Suppose  $y \notin V(F_1)$ . Let  $D = C \cup F_1$ . Let  $F_2$  be a  $(y, D)$ -fan of size three. Let  $\text{end}(F_2) = \{y_1, y_2, y_3\}$ . If  $\text{end}(F_2) \cap (F_1 - \{x_1, x_2, x_3\}) \neq \emptyset$ , then the lemma follows by Claim 1. So we may assume  $\text{end}(F_2) \subset V(C)$ .

*Claim 3.* If  $\{y_1, y_2, y_3\} \subset C^+[x_i, x_{i+1}]$  (If  $i = 3$ , we consider  $x_4 = x_1$ ), then the lemma follows.

*Proof.* We may assume  $y_1, y_2, y_3 \in C^+[x_1, x_2]$  and  $y_1, y_2$  and  $y_3$  appear in this order around  $C^+$ . Then

$$C^+[x_3, y_1] \cdot F_2[y_1, y_2] \cdot C^+[y_2, x_2] \cdot F_1[x_2, x_3]$$

or

$$C^+[x_1, y_2] \cdot F_2[y_2, y_3] \cdot C^+[y_3, x_3] \cdot F_1[x_3, x_1]$$

has the desired properties.

By Claims 1, 2, 3, the only possible case in which the lemma would not hold is  $\{x_1, x_2, x_3\} = \{y_1, y_2, y_3\}$ . We may assume  $x_i = y_i$  ( $i = 1, 2, 3$ ). Let  $D' = D \cup F_2$ . Since  $C$  is a longest cycle,  $C^+(x_1, x_2) \neq \emptyset$ . Since  $G$  is 3-connected, there exists a path  $P$  joining  $C^+(x_1, x_2)$  and  $D' - C^+[x_1, x_2]$  in  $G - \{x_1, x_2\}$ . Let  $\text{end}(P) = \{u, v\}$ ,  $u \in C^+(x_1, x_2)$  and  $v \in D' - C^+[x_1, x_2]$ . If  $v \in V(F_1) \cup V(F_2)$ , then the lemma follows by Claim 2. So we may assume  $v \in C^+(x_2, x_3)$ . Then  $F_1, F_2, C^+[x_1, x_2]$  and  $P[u, v] \cdot C^+[v, x_3]$  form a subdivision of  $K_{3,3}$ . This contradicts the planarity of  $G$ . Therefore, the lemma follows. ■

**Proof of Lemma 2.** Let  $C_0$  be a longest cycle which contains  $x$  and  $y$ . Then  $|C_0| \geq |C|$ . If  $G$  is hamiltonian, then  $C_0$  is a hamiltonian cycle, and  $|C_0| \geq 4$ . Hence the result follows. Therefore, we may assume  $G$  is not hamiltonian, and  $|G| \geq 7$ . By Lemma 1,  $|C_0| \geq \frac{1}{2} \cdot 7 + 2 \geq 5$ . So  $|C_0| \geq \frac{1}{2}|C_0| + 2 \geq \frac{1}{2}|C| + 2$ . Hence we may assume  $z \notin C_0$ . Consider a  $(z, C_0)$ -fan  $F_1$ . Let  $\text{end}(F_1) = \{z_1, z_2, z_3\}$ . We may assume that  $z_1, z_2, z_3$  appear in this order around  $C^+$ . We consider three cases.

Case 1.  $\text{end}(F_1) \subset C_0^+[x, y]$  or  $\text{end}(F_1) \subset C_0^+[y, x]$ .

We may assume  $\{z_1, z_2, z_3\} \subset C_0^+[x, y]$ . Then one of the two cycles  $C_0^+[z_2, z_1] \cdot F_1[z_1, z_2]$  and  $C_0^+[z_3, z_2] \cdot F_1[z_2, z_3]$  has the desired properties.

Case 2. One of  $\text{end}(F_1)$  lies on  $C_0^+(y, x)$  and the other two lie on  $C_0^+(x, y)$ .

We may assume  $z_1, z_2 \in C_0^+(x, y)$  and  $z_3 \in C_0^+(y, x)$ . Let  $C_1 = C_0^+[z_2, z_1] \cdot F_1[z_1, z_2]$ . Then  $C_0 - C_1 = C_0^+(z_1, z_2)$ . Let  $D = C_0 \cup F_1$ . By Theorem B, there exists an  $(x, D - C_0^+(z_3, z_1))$ -fan  $F_2$  of size three, such that  $z_1, z_3 \in \text{end}(F_2)$ . Let  $\text{end}(F_2) = \{z_1, z_3, a\}$ . If  $a \in F_1[z, z_1]$  or  $a \in F_1[z, z_2]$ , let

$$C_2 = C_0^+[z_1, z_3] \cdot F_1[z_3, a] \cdot F_2[a, z_1].$$

If  $a \in F_1[z, z_3]$ , let

$$C_2 = C_0^+[z_1, z_3] \cdot F_2[z_3, a] \cdot F_1[a, z_1].$$

If  $a \in C_0^+(z_2, y)$ , let

$$C_2 = C_0^+[a, z_3] \cdot F_1[z_3, z_2] \cdot C_0^-[z_2, z_1] \cdot F_2[z_1, a].$$

If  $a \in C_0^+(y, z_3)$ , let

$$C_2 = C_0^-[a, z_1] \cdot F_1[z_1, z_3] \cdot F_2[z_3, a].$$

Then in either case,  $C_0^+(z_1, z_2) \subset C_2$  and either  $C_1$  and  $C_2$  satisfies the desired properties. So the only remaining case is  $a \in C_0^+(z_1, z_2)$ . Let  $D' = D \cup F_2$ .

Next, consider a  $(y, D' - C_0^+(z_2, z_3))$ -fan  $F_3$  such that  $\{z_2, z_3\} \subset \text{end}(F_3)$ . Let  $\text{end}(F_3) = \{z_2, z_3, b\}$ . If  $b \in (F_1 - \text{end}(F_1)) \cup C_0^+(z_3, z_1)$ , then the lemma follows by the same argument. If  $b \in F_2(x, a) \cup F_2(x, z_1)$ , let

$$C_3 = F_3[b, z_2] \cdot C_0^-[z_2, z_1] \cdot F_1[z_1, z_3] \cdot F_2[z_3, b].$$

If  $b \in F_2(x, z_3)$ , let

$$C_3 = F_3[b, z_3] \cdot F_1[z_3, z_2] \cdot C_0^-[z_2, z_1] \cdot F_2[z_1, b].$$

Then in either case  $C_0^+(z_1, z_2) \subset C_3$  and hence either  $C_1$  or  $C_3$  satisfies the desired properties. So the lemma follows unless  $b \in C_0^+(z_1, z_2)$ . (Possibly  $a = b$ .)

Now we consider the case  $a \in C_0^+(z_1, z_2)$  and  $b \in C_0^+(z_1, z_2)$ . If  $z_1, b, a, z_2$  appear in this order around  $C_0^+$ , let

$$C_4 = F_3[z_3, b] \cdot C_0^+[b, z_2] \cdot F_1[z_2, z_1] \cdot C_0^-[z_1, z_3]$$

and

$$C_5 = F_2[z_3, a] \cdot C_0^-[a, z_1] \cdot F_1[z_1, z_2] \cdot C_0^+[z_2, z_3].$$



If  $z_1, a, b, z_2$  appear in this order around  $C^+$ , let

$$C_4 = F_3[z_2, b] \cdot C_0^- [b, z_3] \cdot F_1[z_3, z_2]$$

and

$$C_5 = F_2[z_1, a] \cdot C_0^+ [a, z_3] \cdot F_1[z_3, z_1].$$

Then in either case we have  $\{x, y, z\} \subset C_4 \cap C_5$ ,  $C_0 \subset C_4 \cup C_5$ , and hence  $|C_4| \geq \frac{1}{2}|C_0| + 2$  or  $|C_5| \geq \frac{1}{2}|C_0| + 2$ . So the lemma follows.

Now, we may assume that  $a = z_2$  or  $b = z_1$ . If  $a = z_2$ , then  $F_1, F_2, F_3$  and  $C_0^- [b, z_1]$  form a subdivision of  $K_{3,3}$ . If  $b = z_1$ , then  $F_1, F_2, F_3$  and  $C_0^+ [a, z_2]$  form a subdivision of  $K_{3,3}$ . Hence both contradicts the planarity of  $G$ . Therefore, the proof in this case is complete.

Case 3.  $|\{x, y\} \cap \text{end}(F_1)| = |C_0^+(x, y) \cap \text{end}(F_1)| = |C_0^+(y, x) \cap \text{end}(F_1)| = 1$ .

We may assume  $z_1 = x, z_2 \in C_0^+(x, y)$  and  $z_3 \in C_0^+(y, x)$ . Then either

$$\begin{aligned} C_6 &= F_1[z_1, z_2] \cdot C_0^+ [z_2, z_1], \quad \text{or} \\ C_7 &= F_1[z_1, z_3] \cdot C_0^- [z_3, z_1] \end{aligned}$$

satisfies the desired properties.

Therefore, in each case,  $G$  has a cycle through  $x, y$  and  $z$  of length at least  $\frac{1}{2}|C_0| + 2$ .

■

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