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## An inequality between the diameter and the inverse dual degree of a tree

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## Abstract

Let R(T), D(T) be respectively the radius and diameter of a nontrivial tree T and  $I(T) = \sum_{u \in V(T)} 1/\overline{d}(u)$  be the inverse dual degree, where  $\overline{d}(u) = (\sum_{v \in N(u)} d(v))/d(u)$  for each  $u \in V(T)$ . In this note we prove that

$$I(T) \ge \begin{cases} R(T) + 1/3, & \text{if } D(T) \text{ is odd} \\ R(T) + 5/6, & \text{if } D(T) \text{ is even} \end{cases}$$

with equality if and only if T is a path of at least 4 vertices. This inequality strengthens a conjecture of Graffiti.

Let G = (V(G), E(G)) be a simple, connected graph. The distance d(u, v) between two vertices u, v of G is the minimal length of a path from u to v in G. The diameter D(G) of G is the largest distance between any two vertices of G. The radius R(G) of G is  $\min_{u \in V(G)} \max_{v \in V(G)} d(u, v)$ . If  $S \subseteq V(G)$  and  $u \in V(G) \setminus S$ , then we denote  $d(u, S) = \min_{v \in S} d(u, v)$ . The neighbours of  $u \in V(G)$  are vertices adjacent to u in G and the neighbourhood N(u) of u in G is the set of neighbours of u. Since G is simple, the degree of u is d(u) = |N(u)|. The dual degree of u and the inverse dual degree of G are respectively  $\overline{d}(u) = (\sum_{v \in N(u)} d(v))/d(u)$  and  $I(G) = \sum_{u \in V(G)} 1/\overline{d}(u)$ [2]. When ambiguity arises we use  $d_G(u), \overline{d}_G(u)$ , etc., to emphasize that the underlying graph is G.

The main purpose of this note is to prove an inequality between D(T) and I(T). As a consequence we get an inequality involving R(T) and I(T) which strengthens the following conjecture (see [1, 3, 4, 5] for results relating to Graffiti conjectures).

**Graffiti Conjecture 577** For any (nontrivial) tree  $T, I(T) \ge R(T)$ . (There are examples of G which are not trees such that I(G) < R(G).)

In the following we suppose T is a (nontrivial) tree and  $P = v_0 v_1 \dots v_D$  is a path of maximal length in T, where D = D(T). Then  $d(v_0) = d(v_D) = 1$ . Let

$$a = a(P) = |\{v \in V(T) \setminus V(P) : d(v) \ge 2\}|,$$
  

$$b = b(P) = \begin{cases} |\{i : d(v_i) \ge 3, 2 \le i \le D - 2\}|, & \text{if } D \ge 4, \\ 0, & \text{otherwise} \end{cases}$$
  
and  

$$c = c(P) = \begin{cases} |\{i : d(v_i) \ge 3, i = 1, D - 1\}|, & \text{if } D \ge 2, \\ 0, & \text{if } D = 1. \end{cases}$$

**Theorem**  $I(T) \ge D(T)/2 + a/3 + b/10 + c/12 + 5/6.$ 

A caterpillar is a tree with the property that the removal of all degree-one vertices yields a path, called the spine. Note that if T is a caterpillar, then  $v_1 \dots v_{D-1}$  is the spine. To prove the theorem we need the following lemmas.

**Lemma 1** Suppose T is not a caterpillar (so in particular  $D(T) \ge 4$ ) and u is a vertex not in P such that  $d(u) \ge 2$  and d(u, V(P)) is as large as possible. Let T' be the subtree obtained from T by deleting all degree-one vertices adjacent to u. Then D(T) = D(T') and  $I(T) \ge I(T') + 1/3$ .

**Proof** We first note that all but one neighbours of u have degree one, for otherwise there would be a neighbour w of u not in P with  $d(w) \ge 2$  and d(w, V(P)) > d(u, V(P)), violating the choice of u. Suppose  $N(u) = \{u_1, \ldots, u_m, v\}$  where  $d(u_i) = 1$ , for  $i \in \{1, \ldots, m\}$  and d(v) = r. Then

$$D(T) = D(T')$$
. Let  $\sigma = \sum_{x \in N(v) \setminus \{u\}} d(x)$ . Since  $\sigma + 1 \ge r$ , we have  
 $I(T) - I(T')$   
 $= \sum_{x \in N(v) \setminus \{u\}} d(x) + (-1) = -1$   $+ (-1) + (-1) = -1$ 

$$= \sum_{i=1}^{m} \frac{1}{\bar{d}_{T}(u_{i})} + \left(\frac{1}{\bar{d}_{T}(u)} - \frac{1}{\bar{d}_{T'}(u)}\right) + \left(\frac{1}{\bar{d}_{T}(v)} - \frac{1}{\bar{d}_{T'}(v)}\right) \\ = \frac{m}{m+1} + \left(\frac{m+1}{m+r} - \frac{1}{r}\right) + \left(\frac{r}{m+\sigma+1} - \frac{r}{\sigma+1}\right) \\ = 1 + \frac{1}{m+r} - \frac{1}{r} - \frac{1}{m+1} + \left(\frac{m}{m+r} + \frac{r}{m+\sigma+1} - \frac{r}{\sigma+1}\right) \\ \ge 1 + \frac{1}{m+r} - \frac{1}{r} - \frac{1}{m+1}.$$

Note that 1/(m+x) - 1/x is an increasing function of x and  $r \ge 2, m \ge 1$ . We have from the inequality above that

$$I(T) - I(T') \geq 1 + \frac{1}{m+2} - \frac{1}{2} - \frac{1}{m+1} = \frac{1}{2} - \frac{1}{(m+1)(m+2)} \geq \frac{1}{3}.$$
 Q.E.D.

**Lemma 2** Suppose T is a caterpillar but not a path and  $D = D(T) \ge 4$ . If  $d(v_1) \ge 3$  (respectively  $d(v_{D-1}) \ge 3$ ) and let T' be the subtree obtained from T by deleting all degree-one vertices adjacent to  $v_1$  (respectively  $v_{D-1}$ ) excepting  $v_0$  (respectively  $v_D$ ). Then D(T) = D(T') and  $I(T) \ge I(T') + 1/12$ .

**Proof** Suppose  $d(v_1) = m + 2 \ge 3$ ,  $d(v_2) = r$ ,  $d(v_3) = s$ . Then  $r, s \ge 2$  and D(T) = D(T'). We have

$$\begin{split} &I(T) - I(T') \\ &= \sum_{i=1}^{m} \frac{1}{\bar{d}_{T}(u_{i})} + \sum_{i=0}^{2} \left(\frac{1}{\bar{d}_{T}(v_{i})} - \frac{1}{\bar{d}_{T'}(v_{i})}\right) \\ &= \frac{m}{m+2} + \left(\frac{1}{m+2} - \frac{1}{2}\right) + \left(\frac{m+2}{m+r+1} - \frac{2}{r+1}\right) + \left(\frac{r}{m+r+s} - \frac{r}{r+s}\right) \\ &= m\left[\left(\frac{1}{2(m+2)} - \frac{1}{(r+1)(m+r+1)}\right) + r\left(\frac{1}{(r+1)(m+r+1)} - \frac{1}{(r+s)(m+r+s)}\right)\right] \\ &\geq m\left(\frac{1}{2(m+2)} - \frac{1}{(r+1)(m+r+1)}\right) \\ &\geq m\left(\frac{1}{2(m+2)} - \frac{1}{3(m+3)}\right) \\ &= \frac{m(m+5)}{6(m+2)(m+3)} \\ &\geq \frac{1}{12}. \end{split}$$

Q.E.D.

**Lemma 3** Suppose T is a caterpillar but not a path and  $D = D(T) \ge 4$ . If  $d(v_1) = d(v_{D-1}) = 2$ and let T' be the subtree obtained by deleting all degree-one neighbours of  $v_{\alpha}$ , where  $v_{\alpha}$  is the vertex nearest to one terminal vertex of P such that  $d(v_i) \ge 3$ . Then D(T) = D(T') and  $I(T) \ge I(T') + 1/10$ . **Proof** Without loss of generality we may assume  $\alpha \leq \lfloor \frac{D}{2} \rfloor$ . Let  $u_1, \ldots, u_m$  be all the degree-one neighbours of  $v_{\alpha}$ . We have  $d(v_{\alpha-1}) = 2$ . Let  $d(v_{\alpha-2}) = r$   $(r = 1 \text{ if } \alpha = 2 \text{ and } r = 2 \text{ otherwise})$ ,  $d(v_{\alpha+1}) = s, d(v_{\alpha+2}) = t$ . Clearly we have D(T) = D(T'). If  $D \geq 5$ , then  $t \geq 2$ , hence we have

$$\begin{split} I(T) &- I(T') \\ &= \sum_{i=1}^{m} \frac{1}{\overline{d}_{T}(u_{i})} + \sum_{i=\alpha-1}^{\alpha+1} \left( \frac{1}{\overline{d}_{T}(v_{i})} - \frac{1}{\overline{d}_{T'}(v_{i})} \right) \\ &= \frac{m}{m+2} + \left( \frac{2}{m+r+2} - \frac{2}{r+2} \right) + \left( \frac{m+2}{m+s+2} - \frac{2}{s+2} \right) \\ &+ \left( \frac{s}{m+s+t} - \frac{s}{s+t} \right) \\ &= \frac{m}{m+2} + \left( \frac{2}{m+3} - \frac{2}{3} \right) + \left( \frac{m+2}{m+s+2} - \frac{2}{s+2} \right) \\ &+ \left( \frac{s}{m+s+2} - \frac{s}{s+2} \right) \\ &= \frac{m(m+5)}{3(m+2)(m+3)} \\ &\geq \frac{1}{6} \geq \frac{1}{10}. \end{split}$$

If D = 4, then a straightforward calculation shows that

$$I(T) - I(T') = \frac{1}{6} + \frac{4}{m+3} - \frac{2}{m+2} - \frac{2}{m+4} = \frac{1}{6} - \frac{4}{(m+2)(m+3)(m+4)} \ge \frac{1}{10}.$$

Q.E.D.

Now let us prove the main theorem. If  $T = P_n$ , the path with n vertices, then

$$I(P_n) - D(P_n)/2 = \begin{cases} 3/2, & n = 2\\ 1, & n = 3\\ 5/6, & n \ge 4 \end{cases}$$

If D(T) = 2, then T is a star with a = b = 0, c = 1 and  $I(T) - D(T)/2 = 1 \ge c/12 + 5/6$ . If  $D(T) = 3, T \ne P_4$ , then a = b = 0 and T has exactly two vertices with degree  $\ge 2$ . Suppose the degrees of them are l + 1, m + 1. Then max $\{l, m\} \ge 2$  and

$$\begin{split} &I(T) - D(T)/2 \\ &= \frac{\ell + m + 2}{\ell + m + 1} + \frac{l}{\ell + 1} + \frac{m}{m + 1} - \frac{3}{2} \\ &= \frac{1}{\ell + m + 1} - \frac{1}{\ell + 1} - \frac{1}{m + 1} + \frac{3}{2} \\ &\geq \frac{c}{12} + \frac{5}{6}. \end{split}$$

In the following we suppose T is not a path and  $D = D(T) \ge 4$ . If T is not a caterpillar, let u be the vertex not in P such that  $d(u) \ge 2$  and d(u, V(P)) is as large as possible. Then all but one neighbours of u have degree one. Removing from T all the degree-one neighbours of *u* we get a subtree  $T_1$  with  $D(T) = D(T_1), I(T) \ge I(T_1) + 1/3$ , according to Lemma 1. If  $T_1$  is not a caterpillar, then repeat this procedure until a caterpillar is obtained. It is clear that after *a* steps we get a sequence  $T = T_0, T_1, \ldots, T_a$  such that each  $T_{i+1}$  is a subtree of  $T_i$  and  $D(T_i) = D(T_{i+1})$  and  $I(T_i) \ge I(T_{i+1}) + 1/3$ . So we have  $D(T) = D(T_a)$  and  $I(T) \ge I(T_a) + a/3$ .

If  $d(v_1) \ge 3$  in  $T_a$ , then delete all the degree-one neighbours of  $v_1$  except  $v_0$ . We get  $T_{a+1}$ with the same diameter as T such that  $I(T_a) \ge I(T_{a+1}) + 1/12$ , according to Lemma 2. If  $d(v_{D-1}) \ge 3$ , we do the same thing. In this way c subtrees are added to the sequence above and we get  $T = T_0, T_1, \ldots, T_a, \ldots, T_{a+c}$  with  $D(T) = D(T_{a+c})$  and  $I(T) = I(T_{a+c}) + a/3 + c/12$ .

Now we have  $d_{T_{a+c}}(v_1) = d_{T_{a+c}}(v_{D-1}) = 2$  and  $d_{T_{a+c}}(v_i) = d(v_i)$ ,  $i \notin \{1, D-1\}$ . If  $T_{a+c}$  is not a path, then according to Lemma 3 we can delete all degree-one neighbours of some  $v_{\alpha}$  and obtain a subtree  $T_{a+c+1}$  with  $I(T_{a+c}) \ge I(T_{a+c+1}) + 1/10$ . Repeat the procedure until we obtain a path P. When the process stops we get a sequence  $T = T_0, T_1, \ldots, T_a, \ldots, T_{a+c}, \ldots, T_{a+c+b} = P$  with  $I(T) \ge I(T_{a+c}) + a/3 + c/12 \ge I(P) + a/3 + b/10 + c/12$ . Since I(P) = D(P)/2 + 5/6, as we have just proved it for paths, and since D(P) = D(T), we get  $I(T) \ge D(T)/2 + a/3 + b/10 + c/12 + 5/6$ . This completes the proof.

Note that  $R(T) = \lceil D(T)/2 \rceil$  for any tree T and a, b, c are non-negative integers. Hence we have the following corollary.

**Corollary** For any (nontrivial) tree T

$$I(T) \ge \begin{cases} R(T) + 1/3, & \text{if } D(T) \text{ is odd} \\ R(T) + 5/6, & \text{if } D(T) \text{ is even}, \end{cases}$$

with equality if and only if T is a path of at least four vertices.

This corollary tells us that I(T) - R(T) is bounded below. We point out that it is unbounded above. In fact, for the full binary tree T of height  $h \ge 3$  we have  $I(T) - R(T) = 2^{h+2}/5 - h - 1/4$ , which can be arbitrarily large as h tends to infinity.

## References

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