# A semantically-based formatting discipline for Pascal 

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## THE UNIVERSITY OF WOLLONGONG

## A SEMANTICALLY-BASED FORMATTING DISCIPLINE FOR PASCAL

by

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## DEPARTMENT OF COMPUTING SCIENCE

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# A Semantically-based Formatting Discipline for Pascal 

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## ABSTAACT

The abstract (or semantic) syntax of the Pascal language is identified, and a linear representation for the trees so formed within the iramework of the concrete syntax for that language is imposed. The indentation scheme so formed, augmented with a small number of pragmatic considerations, is compared with several previously proposed formatting schemes for Pascal and an example of the use of this new method is given.

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## AUSTRALIA


#### Abstract

SUMMARY The abstract (or semantic) syntax of the Pascal language is identified, and a Ilnear representation for the trees so formed within the framework of the concrete syntax for that language is imposed. The indentation scheme so formed. augmented with a small number of pragmatic considerations. is compared with several previously proposed formatting schemes for Pascal and an example of the use of this new method is given.


KEY WORDS Programming languages Formatted languages Program readability Pascal Abstract syntax

## INTRODUCTION

Recognition that a program is not just an example of person-machine communication but also, and very importantly, one of person-person communication, has developed ever since the appearance of the first "high-level" languages. Such communications include those between program writer and program maintainer during the life cycle of an item of software, and those between members of a team involved in either development or maintenance.

[^1]Methodological and associated language developments such as constructs supporting structured programming ${ }^{1}$ and data abstraction ${ }^{2}$ take cognizance of this tact. Nevertneless. their mere presence in a program. considered as a string of symbols written line by line does not immediately convey to the reader an understanding of its semantic structure.

This paper attempts to provide a remedy to this situation with empnasis on Pascal ${ }^{3}$. by proposing a layout or formatting discipline to clearly display such structures. This layout can be achieved either by formatting or "pretty printing" programs or by individuals when preparing text in the absence of such an aid. in which case. the term "discipline" decomes particularly meaningful. The choice of Pascal is motivated by its widespread use in teaching environments. where
(a) it is used extensively to develop and teach algorithms to students.
(b) in Doth reading and writing programs, students need all the help they can get. in which case, well-defined and uniform layout rules should be ot considerable Denefit.

Of course. such Denefits can also de seen to apply generally.

## PROGRAM STRUCTURE

The semantic structure of a program i.e. the structure in terms of which its meaning is deduced. can be said to be its abstract syntax ${ }^{2}$. This may be represented by a tree structure with "operators" as nodes and "operands" as sub-trees. For example, the Pascal expression

$$
x+y
$$

nas concrete syntax tree

derived according to the syntax rules of the language. but has the abstract tree

conveying the meaning of the addition of $x$ and $y$. Sucn treatment can of course be given to more "interesting" constructs. For exampie.

## while Condition do Statement

can be said to have the abstract or semantic tree

in which the "while do" operator has two operands "Condition" and "Statement" with the obvious meaning.

This capturing of the meaning of a program is fortunate. because the environment of program presentation, Dy lines on a page, admits to the clear and easy representation of tree structures Dy relative indentation. For example, the tree


The methodology for our treatment of Pascal. then. shall be to recognize a suitabie adstract syntax. and to impose that structure on the concrete syntax as manifested in a particular program by the use of indentation. In detail, the various headers and keywords of the concrete syntax shall take the place of the operators of the abstract syntax.

## PASCAL PROGRAM STRUCTURE

Siblings in the abstract tree are in the same way equivalently related under the influence of their associated operator. In the body of a Pascal program or procedure or function, this relationship manifests itself between

- label declarations
- constant declaratıons
- type deciaratıons
- variable declarations
- procedure and function declarations
- the statements of the body
suggesting the following indentation scheme

```
neading
labels
constants
types
variables
functions
procedures
begin
statement
.
end.
```

Note that
(a) Function declarations are placed prior to procedures as a disciplined aid to program reading. "Forward" declarations can be used to overcome difficulties witn the definition order.
(b) There is no indentation brought about by the presence of begin and end, which are meaningiul only as (syntactic) delimiters that are chosen to group all the statements as a node in this level of the tree. Supposing that such a grouping was at all meaningful, then indentation would have been appropriate. However in our view, each statement should in a sense have direct access to the declarations.
(C) When a program or procedure or function contains sub-procedures or suofunctions, their bodies will be indented, so that only the user accessible part, the neading, will be available at the outer i.e. interesting. level of indentation e.g.

```
procedure first:
declarations
procedure second:
    declarations
    begin
    Statements for second
    - N
    end:
begin
Statements for first
end:
```

so that reading the "Statements for first" to discover the meaning of accessed names, the reader has only to look along the current level of indentation. and only those declarations at that level will be worthy of interest, while those inside second. which are of no concern, are clearly moved out of the line of sight by indentation.

Finally. it should be noted that a program would conceivably appear as

## program

declarations
begin
Statements
end.
where the indentation of all out one line of the text is clearly redundant. Uniess goto statements and labels are being used (see below). it is of course reasonable to omit the indicated indentation in this particular situation.

## PASCAL DECLARATIONS

Noting that we have already covered procedures and functions, we see that we nave recognized in our semantic tree the Pascal syntax scheme of grouping labeis. con-
stants, type and variable declarations. This is simply because we nold that this is a matter of the sensible organisation and classification of declaratives whose individuai texts are physically insignificant for purposes of comprehension.

Label declarations appear as
label

$$
I_{1} \cdot I_{2} \cdot I_{3} \cdot \cdots I_{n}=
$$

where each $I$ is a valid label. The keyword appears as a line by itself, corresponding to the semantic operator, the operands are on the succeeding tine and are indented. Canonically, they snould be on separate lines i.e.
label
1.
13.

In:
but the brevity of each label makes the previous proposal the more pragmatically sound.

The canonical arrangement is more correctly shown with respect to constant definitons:
const
Name $_{1} \quad=C_{1}:$

Name $_{n} \quad=C_{n}$ :
We further make the pragmatic observation that the " $=$ " delimeters be aligned for clarity, which may easily be achieved by using a <tab> character or characters, just as can be used to effect indentation. In this regard it is felt that all indentations should be of equal width. Capital letters may also be used for the "name" so that constants may de clearly identifiable wherever they may appear. For example

```
const
BLANKS = ' '
MAXHEIGHT =6:
MINWEIGHT = 100:
```

With types, schematically

## type

Name ${ }_{1}=T_{1}:$

$$
\text { Name }_{n}=T_{n}:
$$

a new item of interest is introduced: recora definition, including variant recoras.
Because the expression of a type definition. the $T_{1}$ above. is already indented. there need be no further indentation for record ... end e.g.

```
type
    man = record
        age : integer:
        address : array [1..60] of char
        end:
```

The discussion of variant records shall be treated under the "case" statement delow. In a similar manner the schematic for "var" declarations would be:
var
Name ${ }_{1}: T_{1}:$

$$
\text { Name }_{n} \quad: T_{m}
$$

To clearly distinguish variable names from constants lower case letters only should be used e.g.
var

| eyecolor | : string: |
| :--- | :--- |
| height | : integer: |
| weight | : integer: |

## PASCAL STATEMENTS

We have already shown how the compound statement which is the body of a program. procedure or function is not indented with respect to begin ... end implying that for consistency the structuring operations which can contain compound statements be regarded as $n$-ary operators.

This is shown rather well when contemplating the Pascal iterative constructs. For the "repeat" statement, we may perhaps have semantically

representable by

> repeat Statement ${ }_{1}$ : until Statement $_{n}$ Condition

Note now the keywords representing the semantic operator are not indented. Dut that the operands are. By no great stretch of the imagination, it is possible to conceive of the "repeat until" operator as being "curried" i.e. it is applied to the "Condition". producing a new semantic operator, which can be then applied to Statement,... .Statement $n$. Diagrammatically

which becomes

## repeat

Statement ${ }_{1}$ :

## Statement $n$ until Condition

1.e. the simple initial application of repeat ... until to the Condition deserves no speclal indentation.

Formats for the "while" and "for" statements are then easily derived:
while Condition do
Statement
while Condition do
begin
Statement ${ }_{1}$ :
-
Statement $n$
end
for Index:= Expression, to Expression ${ }_{2}$ do [or downtol
Statement
for Index := Expression to Expression 2 do
begin
Statement ${ }_{1}$ :

Statement ${ }_{n}$
end
Note that there is no reason for the "do" to appear on an individual line.
With regard to branching, the "if" statement is simply

or

ylelding

## if Condition then <br> Statement: <br> Statement end

and perhaps
else
begin
Statement ${ }_{k}$ :

Statement end
with begin ... end optional in the case of a single Statement. Clearly the tree structure has two levels. whereas our indentation scheme only has one. because of the fortuitous placement of the non-indented else.

An important semantic phenomenon for which Pascal provides only in the limited way of the "case" statement (see below) is the mult-way branch, which must be simulated in Pascal by a series of nested "if" statements:

```
if Condition, then
    Statement
else
    if Condition_ then
                            Statement2
    else
```

Therefore, in the case where. say Condition ${ }_{2}$ above does not represent a condition of the state of the machine depending upon "not Condition, " but represents an alternative to "Condition,". Inen we suggest:
if Condition ${ }^{1}$ then
Statement,
else
If Condition 2 then
Statement ${ }_{2}$
else

## .



The "case" statement is semantically

and therefore gives
case Expression of branch,
oranch
end
en
the closing "end" is indented to avoid

```
begin
Statement:
```

case
end
end

Each branch in detail is a list of labels followed by a series of statements


ו.e.

$$
\begin{gathered}
\text { Label }_{1}, \ldots . \text { Label }_{n}: \\
\text { begin }_{1} \\
\\
\text { Statement }_{1}:
\end{gathered}
$$

## Statement $m$ end

An "others" or "default" branch that some versions of Pascal provide may de treated as one of these e.g.
others:
begin
Statement ${ }_{1}$ :
Statement
end
For example

```
    case day of
    monday, tuesday, wednesday. thursday
                                    begin
                                    rise:
                    work:
                sleep:
                end:
    friday:
                begin
                rise:
                work:
                    drink:
                sleep:
                    end:
    others:
        begin
        rise:
        sleep:
        end
    end:
The format of the variant record declaration may be derived from the format of the "case" statement. The overall structure is
```


## case tag_field type_identifier of

``` Variant \({ }_{7}\) :
```


## Variant $_{n}$

(note the absence of a terminating end). Each variant appears likewise as

$$
\begin{gathered}
\text { label } 1 \text { body }
\end{gathered}
$$

where body is structured as the declaration of procedure or function formal parameters:

$$
\text { (Name } \left._{1}: \text { Type }_{1}: \ldots: \text { Name }_{n}: \text { Type }_{n}\right)
$$

The potential problem of exceeding the right margin is discussed below in the context of parameter declarations.

As an example:

```
type
planet = (venus, mars, pluto):
alien = record
            age :integer:
            weight : real
            case origin : planet of
                        venus,mars:
                            (arms. heads.legs : Integer:
                                    married : boolean):
                                    pluto:
                            (wheels,gears,cogs:integer)
                            end:
```

The inoentation of labels can be extended to the destination of a "goto" e.g.

## Statement

99:
Statement

## goto 99

We are theoretically guaranteed to have the space to make the label outstanding (the Dody of an entire program may need to be indented in this case).

Finally, of structured statements, is the "with" statement, which is obviously

```
with Record do
    begin
    Statement \({ }_{1}\) :
```

    Statement \(n\)
    end

The remaining statements are the "goto". "assignment", and "procedure call" which are not part of this discussion because they do not involve control structure semantic operators.

## MISCELLANY

A sound pragmatic point concerns the lines which exceed some extreme right margin. determined by paper. screen or physical device width. which occur often in the context of procedure and function parameter specifications and calls, and whose frequency increases as the left margin is indented further.

For procedure and function parameters. the following scheme is proposed. instead ot an norizontal decomposition, we propose

```
procedure P (formal-parameter
    formal-parameter }n\mathrm{ ):
```

lor a procedure neading, and

> function F ( formal-parameter
formal-parameter ${ }_{n}$ )
$\therefore$ function-type:
for a function heading. and
PF ( actual-parameter

$$
\text { actual-parameter } n \text { ): }
$$

for a procedure or function call. This is done purely for the sake of aesthetics. For exampie

| procedure readyfile filename | : namestr: |
| ---: | :--- |
|  | var numberoffields |
|  | : integer: |
|  | : field |

Should an entity (typically, part of a long expression. or pernaps part of a formai parameter specification) exceed the right margin. a break is simply made at the last blank part of the right margin and contınuation proceeds from the current indent position on the next line e.g.

$$
\begin{aligned}
\text { verylong := } & (\text { longer * verylong ) ) } \\
& \text { summationofvariables: }
\end{aligned}
$$

Note now the continued expression lines align from the left. and not the position of the control line.

Finally, we remark that the advantages of horizontal spacing, particularly detween groups of declarations, cannot be overestimated. and that similarly mandatory snould De a comment explaining the nature of a procedure, function or program after its neading.

## PREVIOUSLY PROPOSED SCHEMES

The merits of our proposal, apart from its semantical foundation and deauty, can de demonstrated by discussing some previously proposed schemes.

Chronologically, the first published formatting scheme for Pascal is that which could De deduced from the pages of Wirth ${ }^{4}$. which is that presumably called the "Classical" style $^{5}$. While the examples presented are rather inconsistent in indentation, nonetheless the following trends are observed:
(a) the general structure of a program/procedure/function is

```
neading
    declarations
begin
    statements
end
```

which vaguely corresponds to ours save for the irrelevant prominence of the begin ... end brackets.
(b) the bodies of loops and trailers tend to be indented; if they are single statements they sometimes appear on the same line as the semantic operator e.g.

## if Condition then Statement <br> else Statement

and the indentation of a compound statement body is apparently a consequence ot begin ... end influence e.g.

## If Condition then begin <br> Statement

end
However, we repeat that such rules are deductions, and require clarification and improvement - the currently described piece of work can be thought of as such.

Some concrete proposals can be found in Singer. Hueras and Ledgard ${ }^{6}$. Theır merit is the emphasis on use of vertical spacing to aid clarity, and on the definition of various disciplines regarding interfacing of sub-programs to their environments (such issues are regarded as being outside the scope of the present work). On the other nana
(a) not much is said about the indentation of nested procedure definitions. Is the following forbidden?

```
proceaure A
procedure B
begin
    Statement
end:
begin
    Statement
```

end
(b) while the bodies of loops etc are to be indented, so are the subordinates of a begin ... end which grouping we have shown to de semantically insignificant. always occurring as part of a more meaningful set. The present scheme allows double indentation

# while Condition do begin <br> Statement 

end
which can only lead to a quicker than necessary confrontation with the right margin. Similarly superfluous is the required indentation of the detail of a record definition between record and end. The detail of the indentation of the bodies oi a "case" statement with respect to the case labels (and labels in general) is unspecified. Finally, the then and else branches of the "if" are given yet more indentation e.g.

## If Condition <br> then

begin
Statement
end
else
begin
Statement
end
which we suggest demonstrates the lack of a formal correspondence between this scheme and the abstract syntax of its object programs.

A further set of rules in Peterson ${ }^{7}$ which does not address the general methodological questions we felt the latter did so well echoes its mistakes with regard to indentation. Notable is a lack of specification of a number of concepts and relationships such as nested procedure declarations (which is an absolutely vital issue in the environment of top-down development and the layout of some statements and labels. Present again is the idea that begin ... end is semantically significant ( which significance is represented by indentations). Even more obscure is the way in which multiple levels of
indentation are required in the bodies of the control constants, tor example
If Condition then begin

Statement

## end <br> else begin

Statement
and end $\begin{gathered}\text { while Condition } \\ \text { do begin } \\ \text { Statement } \\ \text { end }\end{gathered}$
which is faulty in that
(a) there is an excess of syntactic detail and complex indentation to introduce some essentially simple concepts.
(b) indentation to the right proceeds excessively quickly.
(C) the right-justification of while and do must be hard to effect. similarly, the indentations are not uniform, but depend upon the relevant semantic operator.

An interesting scheme ${ }^{5}$ addresses the issue of indentation in addition to the irritation caused to the Pascal programmer by having to remember the trivial details of placement of semicolons and the need to insert begin ... end around compound statements as Dodies of such as while's and if's. The latter is solved by the elegant expedient of making begin and end part of every structured construct. ana. because Pascal allows an empty statement, placing a semicolon after every statement (i.e. betore "end" - we see that a semicolon can never occur before an else, thougn). For exampie, we write

## If Condition then begin <br> Statement <br> end else begin <br> Statement <br> end

where the Statement can be a simple or multiple statement without fear. The scneme admittedly addresses only part of the layout problem, and could be incorporated into our iramework as an alternative to our particular formulations in this regard. It is after much consideration that we choose not to, because
(a) the fate of the outer begin ... end for the body of a program/procedure/function is not simply and uniformly accounted for (it causes an indent of its own in the scheme referred to)
(b) we are in a sense changing the language, hypothesising a new concrete syntax along the lines of

## If Condition then <br> Statement <br> else <br> Statement <br> - <br> fi

and implementing it in terms of combinations of elements of the Pascal syntax: we choose not to impose this burden on the programmer believing our scneme aids layout just as effectively and in a more familiar and natural way.

The work of Singer et $\mathrm{al}^{6}$ and Peterson ${ }^{7}$ is further discussed by Mohilner ${ }^{8}$ with respect to the "right margin" problem. We feel that our solution to this problem is on firmer foundations. Oppen ${ }^{9}$ discusses pretty printing from an implementation viewpoint rather than discussing a basic formatting policy for particular styles of language.

Our final example is an analysis by Rose and Weish ${ }^{10}$ of a version of Pascal with inouilt formatting rules as part of the syntax. avoiding the need for begin ... end
delimiters and semicolons and the end terminator (of case and record entities). We see the basic principles of our design appearing. but of course without the need to consider the placement of begin and end (the possibility of the omission of which. as this work demonstrates, proves their uselessness as a semantically meaningful construct. Details with which we quibble are
(a) the layout of the "if" statement as
if Condition then

Statement
else
Statement
there being no reason for the placing of the "then" on a single line.
(D) sometımes we then see

> If Condition then Statement else Statement ${ }_{1}$ :   Statement $_{n}$
which masks the "then Statement" and similarly in the case of an "else Statement".
(C) multiple statements per line. which mars clarity
(d) similarly, (case) labels and statements on the same line.

In summary these proposals are sound, particularly with respect to the way in whicn they agree with ours about the nested structure of declaratıons. but we claım the reiative merit of our work is in the accommodation of the actual Pascal syntax, and in the strictness of our approach with respect to the vertical separation of text.

## CONCLUSION

Ease of understanding of a program is facilitated by a clear exposition of its semantic tree structure, which can be achieved by indentation. A formatting discipline ror

Pascal based upon indentation to reflect the semantic structure of programs has deen developed. A comparison with other schemes of varying quality has Deen attempted. the results of which comparison by themselves command the proposed scheme. The scheme is simple, straightforward and semantically elegant.

## EXAMPLE

program example (input, output:
( example -- reads a list of not greater than 1000 numbers and outputs them in ascending order)
const
LIMIT = 1000: ( upper bound on number of items read)
type minrange $=1 .$. LIMIT; (index type for table $)$
var
numbersread : 0..LIMIT: ( counter for numbers read) table : array [minrange] of integer; ( number storage ) 1 : minrange: ( loop index )
proceaure smallfrom (base : minrange):
( smallfrom -- select the smallest number in the range "base" to "numbers-read")
var
min : minrange; ( index of smallest number found)
function indexsmallfrom (base : minrange) : integer:
( indexsmallfrom -- integer function to give the position of the smallest number starting at position "base"\}
var
index : minrange: (index of smallest number found so far )
Degin
if base $=$ numbersread then
indexsmallfrom:= base
eise
Degin
index := indexsmallfrom (base + 1):
if table [base] < table [index] then
indexsmallfrom:= base
else
indexsmallfrom:= index
end
ena:
Degin
mın := indexsmallfrom (base):
writeln (table (minj):
table [min] := table [base]
end:
Degin
numbersread := 0 :
while not eof do
begin
numbersread := numbersread + 1:
readin (table [numbersread])
ena:
for $1:=1$ to numbersread do
smallifrom (i)
ena.

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