# Quantitative LEED analysis using a simultaneous optimisation algorithm.

# M. Blanco-Rey<sup>1,3</sup>, K. Heinz<sup>2</sup>, P.L. de Andres<sup>1</sup>

- <sup>1</sup> Instituto de Ciencia de Materiales (CSIC), Cantoblanco, 28049 Madrid, Spain
- $^2$  Lehrstuhl für Festkörperphysik, Universität Erlangen-Nürnberg, Staudtstrasse 7, 91058 Erlangen, Germany
- $^3$  Department of Chemistry, University of Cambridge, Lensfield Road, Cambridge CB2 1EW, UK

Abstract. The performance of a combinatorial simultaneous optimisation algorithm (SO) is tested using experimental LEED I(E) data from Cu(100) and Fe<sub>0.57</sub>Al<sub>0.47</sub>(100) surfaces. SO optimises structures taking advantage of the experimental database at two levels: (i) conmensurate subsets of the database with the number of unknown parameters are chosen to find local solutions using Broyden's method and, (ii) these partial structural solutions are used to build a Markov chain over the whole database. This procedure is of global character, the same as simulated annealing or genetic algorithms methods, but displays a very competitive scaling law because after the first iteration candidates are not chosen by a blind/random pick; they are already solutions to the problem with a restricted experimental database.

PACS numbers: 61.14.Hg, 68.03.Hj

Submitted to: JPCM

#### 1. Introduction

Most properties of surfaces (e.g. their reactivity) are strongly influenced by their structure. A number of diffraction techniques have been developed in Surface Science to find this structure with atomic accuracy, such as Low-Energy Electron Diffraction (LEED) [1, 2, 3], Photoelectron Diffraction (PD) [4], Surface X-ray Diffraction (SXRD), [5, 6], Near-Edge X-ray Absorption Fine Structure (NEXAFS), [7] etc. [1, 2, 8, 9]. Among these techniques, LEED plays a prominent role in the field due to similar progress made in experimental and theoretical techniques, and nowadays is a technique that can be found in nearly every surface science laboratory. Experimental intensities for different Bragg beams diffracted from the surface are measured as a function of the energy of the primary beam of electrons; these carry the information on the geometrical position of atoms in the surface.

Diffraction amplitudes are complex-number functions of the parameters defining the structure; these functions may be multivalued and their inverse has to be found inside the physical appropriate Riemann sheet (the existence of several branches is usually referred in the literature as "multiple coincidences"). In principle nothing prohibits to find the inversion algorithm linking the structure and the measured intensities, which would provide us with the geometrical parameters in just one single step. In practice, several technical problems make this a difficult program; we shall only mention the following two difficulties: (1) The impossibility of measuring the diffracted wavefield phase makes the inversion difficult as a great deal of information is carried away with the phase itself and, (2) the strong interaction between the incident wavefield with the surface electrons, i.e. multiple-scattering, makes the direct/inverse functions highly nonlinear, and quite difficult to localize the physically meaningful branches. Therefore, the standard approach over the years has been to minimize a cost function, the so-called R-factor, designed to quantify the agreement/discrepancy between experimental and calculated spectra. This in turn is not free of complications because the R-factor is an scalar function of a vector argument defining all the structural and non-structural parameters needed to define the system, leading to a non-polynomial scenario in the search for the global minimum that only could be solved with complete certainty in an infinite amount of searching time. We should mention here how Prof. John Pendry and his collaborators have led the way to identify and to find solutions to all aforementioned fundamental problems for the LEED technique[1, 10, 11, 12, 13].

In this paper we analyze a new route to overcome the intrinsic non-polynomial (NP) nature of traditional LEED analysis based on the use of an R-factor, in particular Pendry's R-factor[11]. The method is applied to experimental data, and its performance is compared with state-of-the-art solutions proposed in the literature. More precisely, we want to solve the following problem: Let N be the number of parameters defining the structure. Then, the R-factor is a scalar function defined in  $\Re^N$ . In general, this hypersurface displays a complex topography, showing several local minima. If we intend to find out its global minimum by simple brute force search, i.e. evaluating the R-factor

for every point of a grid in parameter space, we find that the number of trial model structures, t, scales exponentially with N. In the case of a grid where the number of values in each of the N dimensions is equal to  $n_p$ ,  $t = n_p^N$  structures should be tried. Because the time needed to show the structure scales exponentially with the number of unknown parameters, optimisation theory classifies this as a non-polynomial or NP-complete problem [13]. In particular, if we carry out an exhaustive search on a continuous interval of length L for each parameter, making sure that the solution is at the worst case at a maximum distance  $\epsilon$  from the true solution, we obtain:

$$t \simeq \left(\frac{L\sqrt{N}}{\epsilon}\right)^N \tag{1}$$

For high N values, this scaling law,  $t \sim N^{N/2}$ , is even worse than the exponential law that appears in the grid case.

We will restrict ourselves to the determination of an ordered surface structure from quantitative LEED. For a single incidence direction, the database is formed by a set of I(E) curves. These are intensity vs. energy curves, one per reflection or "beam". For a given model surface, the accurate evaluation of I(E) requires using a multiple scattering formalism, which is computationally expensive. Improvements in LEED I(E) analysis efficiency made so far have been focused in the development efficient intensity evaluation methods, such as Tensor-LEED (TLEED) [12], and in the development of algorithms that reduce the number of intensity evaluations during the optimisation procedure. The present paper deals with the second subject.

In their pioneering contribution to develop direct methods for LEED, Pendry, Heinz and Oed introduced the idea of reducing the traditional optimization method to a multidimensional root finding. Starting from a linearized version of TLEED for the scattering amplitudes, they have proposed an iterative solution for the related nonlinear system of equations for intensities[13]. This class of methods fits all the datapoints simultaneously, since it involves the minimisation of a vectorial function  $f: \mathbb{R}^N \to \mathbb{R}^N$ . Indeed direct methods based on the holographic approach[14], and the quasi-direct method proposed by the direct fitting of a subset of experimental intensities have been the two most promising routes to find solutions for structural problems that could easily get out of hand as the system becomes more and more complex.

Multidimensional function minimisation techniques can be classified into local and global. Local methods (e.g. steepest descent making downhill moves in the R-factor hypersurface [15, 16, 17]) have the advantage of scaling like the number of unknowns to be found squared  $(N^2)$ , but they have the disadvantage of exploring only the nearest single local minimum. Alternatively, global methods allow uphill moves on the hypersurface to explore several minima by crossing between "valleys". Most of these methods work by constructing Markov chains of structures and are based in the ergodic principle. Some examples are Genetic Algorithms (GA) [18], Random Sampling Algorithms (RSA) [19] and Simulated Annealing (SA)[20]. All of them would find the correct global minimum if given infinite time; however, within a finite time interval

the degree of success varies depending on the particular circumstances and has to be studied independently, often to optimize the parameters in the search. As an example of selection of different parameters, the simulated annealing algorithm is interpreted as a physical adiabatic cooling performed under quasi-equilibrium conditions, and the cooling scheme is choosen to achieve a robust performance [20, 21].

In a previous paper, we have proposed simultaneous optimisation (SO)[22] based on a number of theoretical simulations, and we have shown it can be successful in two highly nonlinear scattering problems: single atom phase shifts retrieval from backscattered intensities, and surface structure optimisation from LEED simulated data. In the present paper, the SO method is described in full detail and its performance is tested against real LEED experimental I(E) data measured for Cu(100) and Fe<sub>0.53</sub>Al<sub>0.47</sub>(100) surfaces. These two systems have also been studied with RSA, allowing us to take a benchmarck on performance against a different method known already to be efficient and robust. The SO algorithm consists of two shells: (i) an inner shell where structures are obtained by solving non-linear systems of equations by applying Broyden's method[23] (at this point only a random subset of the experimental database is used), and (ii) an outer shell, where the structures obtained in the inner shell are validated against the whole experimental database by using Pendry's R-factor.

## 2. The algorithm

Several R-factor prescriptions can be considered for the structural problem[2]; the basic task consists in comparing values for calculated and experimental intensities curves, I(E), but first and/or second-order derivatives of the I(E) curves might be used too. Let us consider a simple and basic R-factor, the least-square correlation factor:

$$R \sim \sum_{\vec{\alpha}} |q^{th}(\vec{x}, \vec{\alpha}) - q^{exp}(\vec{\alpha})|^2 \tag{2}$$

where q is used to indicate that any convenient function of the intensities might be used for the comparison,  $\vec{x} = (x_1, \dots, x_N)$  are the N unknown structural parameters, and  $\vec{\alpha} = (E_i; \vec{g})$  runs over energies and beams. The full experimental database is made of intensities measured for different energies,  $E_i$ , in every accesible beam,  $\vec{g}$ :  $\mathcal{S} = \{I_{\vec{\alpha}}\}$ , containing  $N_D$  numbers.

From quite general statistical arguments, we expect the R-factor hypersurface to have a number of shallow local minima, related to random correlations between theory and experiment, and a few deeper wells due to structural coincidences [20]. Thus, if a downhill method were to be used, an independent minimisation would be needed at every local minimum. Since the number of minima grows exponentially with N, such a multiple launch strategy is an NP-problem.

A most popular R-factor,  $R_P$ , has been defined by Pendry [11]. Intensities are affected by background contributions of temperature effects and electron damping, so R-factors that make use of I(E) derivatives are, in general, more reliable because these contributions are largely cancelled out. In particular,  $R_P$  treats all maxima

and minima at the I(E) equally, since all of them contain useful information about constructive/destructive interference conditions closely related to the structure. In fact, Pendry's R-factor compares the so-called Y-functions, based on logarithmic derivatives, to highlight all the extrema found in the spectra, and  $R_P$  quantifies the agreement between experimental and calculated Y-functions.

Usually, the set of experimental data, S, contains many more datapoints,  $N_D$ , than geometrical parameters in the model, N. As pointed out by Kleinle et~al, R-factors using the information contained in their derivatives need to be calculated on a fine grid of energies where some datapoints are correlated. These authors suggested computing an R-factor that only depends on intensities on a coarse energy gride[24]. From a formal point of view, only a small subset of N independent datapoints  $S_N \in S$  should be needed to uniquely determine the structure. In practice, however, multiple coincident solutions might be found for a reduced subset of the experimental database. A characteristic for the global solution is that it it is independent w.r.t different choices for the subset  $S_N$ . This kind of procedure constitutes a list of structures converging to the "global" one, where the starting points do not need to be chosen at random, but are already partial solutions to the global problem. The SO algorithm has been designed to evolve and validate solutions obtained with  $R_P$  from reduced database subsets.

We find that in most cases, these restricted subsets already contain the relevant topographic features of  $R_P$ , although they might also display features not related to the global hypersurface. Obviously, these solutions constitute far better trial structures than simple random moves on the  $R_P$  hypersurface. However, dealing with experimental or noisy data makes  $R_P \neq 0$  at the global minimum. Therefore, the global solution is not necessarily a solution of every partial subset, and determining the global solution can be more difficult, in particular the algorithm could be trapped oscillating around different local minima. It is possible to damp and minimize these oscillations by using overdetermined systems, because by increasing the size of the system of equations the local topography becomes more and more similar to the global one. We have found that for overdetermined systems of the order of  $N_D = 2N$  or  $N_D = 3N$ , the computing effort at the inner shell remains reasonable, while the efficiency increases considerably.

Having in mind that our goal is to find the global minimum in  $R_P$ , we first define a database subset,  $\tilde{\mathcal{S}} \in \mathcal{S}$ , such that it includes intensities where  $\frac{\partial I(E;\vec{g})}{\partial E} \approx 0$  for at least one of the experimental curves (let us call the cardinal of this subset  $\tilde{N}_D \geq N$ ). This choice of  $\tilde{\mathcal{S}}$  is arbitrary, but justified by the fact that those are the points contributing more to Pendry's  $R_P$ .

A system of  $N_{eq}$  equations and N unknowns is constructed with  $N_{eq}$  intensities chosen at random inside  $\tilde{S}$ . We denote this subset by:  $S_{N_{eq}} \in \tilde{S}$  ( $\tilde{N}_D \geq N_{eq} \geq N$ ). Thus,  $\binom{\tilde{N}_D}{N_{eq}}$  different choices of systems of equations are available. We define the following function:

$$f_i(\vec{x}; E_i) = \sum_{\vec{q}} |Y_{\vec{q}}^{\text{th}}(E_i; \vec{x}) - Y_{\vec{q}}^{\text{exp}}(E_i)| = 0 ; \quad i = 1, \dots, N_{eq}$$
 (3)

where  $(E_1, \ldots, E_{N_{eq}})$  is a set of  $N_{eq}$  different experimental points taken from  $\tilde{S}$  at random. Summation over beams is performed to gather as much information as possible, and because it results in a smoother function than the one corresponding to the  $\vec{g}$  value associated to the extreme in the curve. This is relevant, since derivatives need to be calculated in the inner shell (see Sec. 2.2).

#### 2.1. The outer shell: the Markov chain

The global search starts with a structure chosen at random inside the physical accesible region. These will be validated in the outer shell by an iterative process. The k-th iteration in the outer shell starts by calling the inner shell with a structural candidate,  $\vec{x}^{(k)}$ , and a subset of experimental datapoints,  $S_{N_{eq}}$ , chosen at random in  $\tilde{S}$ . The inner shell returns as output a structure,  $\vec{x}'^k$ . Before computing  $R_P$  for the global database, which is more expensive from a computational point of view than invoking the inner shell, it is worth to check that (i) the new structure,  $\vec{x}'^k$ , is also a solution for a different choice of  $S_{N_{eq}}$ , and (ii) whether the new solution stays within a given convergence radius from the old one.  $N_{PC} \sim N$  inner shell calls are made inside the k-th iteration on the outer shell to reduce the number of times  $R_P$  needs to be computed.

The k+1-th iteration begins with a structure,  $\vec{x}^{(k+1)}$ , incorporating small random modifications over the solution in the k-th iteration. Unlike SA, the modifications are always accepted, which is known in the literature as blind random search iterative improvement, and it is a global search method[25]. Following the RSA idea[19],  $\vec{x}^{(k+1)}$  is obtained by considering a gaussian probability distribution centred at the solution in the k-th iteration, where the width of the gaussian,  $\sigma$ , depends on the value of  $R_P(\vec{x}'^k)$ , computed at the end of the k-th iteration. Therefore, uphill moves are guaranteed and the search space size is reduced for small values of  $R_P$ , where long jumps are not necessary.

The partial problem tackled in the inner shell is highly non-linear; it might happen that the system solver does not find the right solution. However, wrong trial structures coming from the inner shell usually have already a few useful components, making better candidates than a pure random blind choice made by conventional algorithms based in the ergodic theorem. Non-gaussian probability distributions, such as Cauchy-Lorentz, have been used in the literature to generate a Markov chain of structures[21]. In our case, we have seen that the distribution shape is not crucial since the system solver method in the inner shell is already quasi-global, and it overcomes  $R_P$  hypersurface simple barriers.

## 2.2. The inner shell: the system solver

To deal with the problem of solving the highly nonlinear system  $\vec{f} = 0$ , a globally convergent algorithm that combines a multidimensional secant method with a descent strategy is used [26]. The only requirement is that  $\vec{f}$  is a continuous and differentiable function w.r.t.  $\vec{x}$ . The *i*-th iteration move in a conventional multidimensional secant

method,  $\delta \vec{x}^{(i)} = \vec{x}^{(i+1)} - \vec{x}^{(i)}$ , is given by:

$$\tilde{J}^{(i)} \cdot \delta \vec{x}^{(i)} = -\vec{f}^{(i)} 
\tilde{J}^{(i+1)} \cdot \delta \vec{x}^{(i)} = \delta \vec{f}^{(i)}$$
(4)

where  $\delta \vec{f}^{(i)} = \vec{f}^{(i+1)} - \vec{f}^{(i)}$  and  $J_{jk}^{(i)} = \frac{\partial f_j^{(i)}}{\partial x_k^{(i)}}$  are the Jacobian matrix elements. In order to save derivative evaluations, we substitute the exact Jacobian by Broyden's approximation [23],

$$\tilde{J}_B^{(i+1)} = \tilde{J}_B^{(i)} + \frac{(\delta \vec{f}^{(i)} - \tilde{J}_B^{(i)} \cdot \delta \vec{x}^{(i)}) \otimes \delta \vec{x}^{(i)}}{\delta \vec{f}^{(i)} \cdot \delta \vec{f}^{(i)}}$$

$$(5)$$

rather than the exact Jacobian  $\tilde{J}^{(i+1)}$ . The procedure is initialised with the exact Jacobian. When the system  $\vec{f} = 0$  is overdetermined, the linear system of equations Eq. (4) is solved in a least squares sense.

However, this method tends to wander around the parameter space if the starting point is not close enough to the solution. This can be avoided by embedding the procedure in a globally convergent strategy that tries to minimise the scalar function  $|\vec{f}|^2$ . The latter defines a hypersurface that can be seen as a projection of the global  $R_P$  onto the data subspace  $\mathcal{S}'_{N_{eq}}$ . If we use the exact Jacobian  $\tilde{J}$ ,  $\delta \vec{x}^{(i)}$  yields already a descent direction for  $|\vec{f}|^2$ , but the whole step may not make  $|\vec{f}|^2$  smaller. To ensure it, a line search is performed along this vector,

$$\vec{x}^{(i+1)} = \vec{x}^{(i)} + \lambda \delta \vec{x}^{(i)} \tag{6}$$

taking the  $\lambda$  value that yields the smallest value of  $|\vec{f}|^2$ . Sophisticated line searches could be made at this stage [26], but an equispaced sampling along the line is found to be enough for our purposes. Finally, at the end of each iteration, the exiting criterion is checked. We use convergence in the modulus of  $\vec{x}$  as exiting condition:

$$|\vec{x}^{(i+1)} - \vec{x}^{(i)}| < x_{\min}$$
 (7)

This method is non-local, because, even if the line search takes place along a descent direction of  $|\vec{f}|^2$ , a longer step can overcome a hypersurface barrier. If  $\tilde{J}_B^{(i+1)}$  deviates from the descent direction, the procedure can be restarted with the exact Jacobian.

Every solution to Eq. (3) minimises  $|\vec{f}|^2$ . However, Eq. (3) may eventually have multiple solutions, and only one of them coincides with the global minimum of the R-factor. The others correspond to local minima of the R-factor. In fact, the  $|\vec{f}|^2$  hypersurface shares topographic features with the R-factor.

However, that is true only for the ideal case of perfect agreement between theory and experiment at the global minimum, where  $R_P = 0$ . In a realistic scenario, the global minimum has  $R_P > 0$ . This implies that there might exist partial problems that do not posses a solution at the global minimum. Adding more datapoints to the partial problem should filter out oscillations near the R-factor minima and reconcile the partial problem topographies. We define the oversampling factor  $N_{ov}$  for the partial problem such that the systems contain  $N_{eq} = N_{ov} \times N$  equations and N unknowns, and solve

linear Eqs. 4 in a least-squares sense. Although  $N_{ov}$  should be chosen for each specific problem according to the fit quality, the condition  $N_D >> N_{eq}$  must hold in order to keep the partial problem tractable at a low computational cost.

# 2.3. Scaling law estimate

By construction, each Broyden approximate Jacobian evaluation requires 2N intensity evaluations, while the exact Jacobian would need  $N^2$ . The line search makes a constant number of evaluations, and  $R_P$  is evaluated every  $N_{PC} \sim N$  partial problem calls. Thus, the total scaling exponent is  $\beta \approx 2$  before entering the Markov chain. On the other hand, we can expect the outer shell to have a maximum scaling similar to RSA algorithm, which has an exponent  $\beta \approx 2.5$  [19]. Notice that the present estimate considers the number of intensity evaluations only, and not other N-dependent operations present in the algorithm that affect the total computation time, such as matrix inversions. This analysis gives an upper boundary value for the exponent,  $\beta \lesssim 4.5$ . SO yields efficient scaling laws because the structures entering the Markov chain are by construction partial solutions to the global problem, i.e. the outer shell is similar to an importance sampling, rather than to a standard random sampling. In fact, we have already seen in the statistic results from noiseless simulated LEED I(E) data of the Ir(110)- $p(2\times1)$  surface, that SO yields an exponent  $\beta = 4.1 \pm 0.1$  [22]. A comparable value,  $\beta = 3.5 \pm 0.2$ , was found in the atomic phase-shifts retrieval from noiseless backscattered intensities from a single atom[22].

From the arguments above, it is expected that the use of overdetermined nonlinear systems will affect the scaling exponent only slightly, though the total computation effort will increase. This has been observed in the problem of atomic phase shifts retrieval from single atom noisy backscattered intensities[22]. In that example, a convincing success rate could be obtained at low values of  $N_{ov}$ . Adding 5% of gaussian noise in the intensities and applying oversampled SO resulted in a scaling law exponent  $\beta = 4.4 \pm 0.3$ , this value slightly increased to  $\beta = 4.8 \pm 0.4$  when noise was 10%. For both noise values, success rate above 90% could be routinely obtained using  $N_{ov} = 2$  when searching for  $l_{\text{max}} < 5$  phase shifts, and  $N_{ov} = 4$  for  $l_{\text{max}} \ge 5$ .

The main quantity used to describe the efficiency of SO is the computing effort, defined as the expected value

$$\langle t \rangle = \frac{\langle N_{call} \rangle}{N_e} \tag{8}$$

where  $N_{call}$  is the number of individual intensity evaluations and  $N_e$  is the number of energy points per I(E) curve. In those previously published theory-theory benchmarks[22], it was found that for each value of N the probability distribution function (PDF) of a number  $N_{stat}$  of independent searches shows a peak at low t values and is exponentially decreasing. Interestingly, it can be fitted to a one-event Poisson-like distribution:

$$p(t) = w^2 t e^{-wt}. (9)$$

This is the PDF of a single random event taking place in a predetermined time interval between t and t + dt, the event being that SO finds the global minimum. The area under p(t),

$$P(t) = \int_0^t p(t')dt' = 1 - (1 + wt)e^{-wt}$$
(10)

gives the success rate of a search that uses a computing effort t.

#### 3. Results with experimental data

# 3.1. Clean Cu(100) surface

SO performance has been benchmarked using experimental normal-incidence LEED I(E) curves from a clean Cu(100) surface. Pendry's R-factor,  $R_P$ , is used to quantify the agreement between experimental and calculated spectra. The experimental data were taken at a low temperature (90 K) in order to reduce thermal diffuse scattering (details can be found elsewhere [27]). The Sample quality was good (common impurities, such as C, O and S, were below the Auger detection limit) and the I(E) curves were reproducible with  $R_P < 0.02$ .  $R_P$  values among symmetry related beams lie below 0.04, ensuring a correct sample alignment. Previous dynamical analyses of this surface have been reported in Refs. [27] and [28]. It was found that it is necessary to introduce an energy dependent inner potential  $(V_{or}(E))$  in order to accurately reproduce the correct lattice parameter,  $a_0 = 2.55 \text{ Å}$  [28]. A theoretical model for  $V_{or}(E)$  can be parametrised from first principles [29]. When a constant  $V_{or}$  is used, the I(E) fit yields a contracted in-plane lattice parameter value of  $a_0 = 2.53 \text{ Å}$  [27]. This deviation of cents of Å lies beyond the  $R_P$  variance limit due to the systematic error caused by neglecting  $V_{or}(E)$ . To discuss the efficiency of the SO algorithm when fitting geometrical data, for simplicity, we will restrict ourselves to a constant  $V_{or}$ , use the smaller  $a_0 = 2.53$  Å and keep Debye parameters fixed at their optimum values. The database used in this work consists of four beams, namely (10), (11), (20) and (21), measured in the energy range 55-410 eV and producing a total data base width of  $\Delta E = 940$  eV. The maximum energy has been chosen to be smaller than the 500 eV used in the original calculations [27], so as to make the  $V_{or}(E)$  dependence less relevant for the fit. A maximum angular momentum number  $l_{max} = 11$  achieves convergence such that differences in  $R_P$  are smaller than 0.001.

An energy-averaged value of  $V_{or}$  is also to be fitted. Conventional directed algorithms act upon the whole I(E) curve, so that it is usually enough to apply a rigid shift to the I(E) curve to make peaks coincide and obtain a good estimate of  $V_{or}$ . The straight-forward way of introducing  $V_{or}$  in a SO search would be to use it only in the outer shell of the algorithm. In the first iterations, when calculated and experimental I(E) curves are poorly correlated, the estimated value of  $V_{or}$  will be, in general, in disagreement with the actual value. If this incorrect  $V_{or}$  is kept fixed during inner shell iterations, it will prevent the algorithm from approaching the solution. A better strategy consists in considering  $V_{or}$  as an extra unknown in the inner shell, thus

adding another equation to the system. We have used this second procedure in the forecoming calculations.

As a reference, we will use the structure determined full-dynamically by Müller et al [27], with  $a_0 = 2.53 \pm 0.01$  Å for the lattice parameter, and the following interlayer distances (ordered from vacuum to bulk):  $d_1 = 1.765 \pm 0.005 \text{ Å}, d_2 = 1.805 \pm 0.010 \text{ Å},$  $d_3 = 1.80 \pm 0.01 \text{ Å}, d_4 = 1.79 \pm 0.02 \text{ Å}, d_5 = 1.80 \pm 0.03 \text{ Å}, d_6 = 1.79 \pm 0.04 \text{ Å}, and$ bulk interlayer distance  $d_b = 1.79 \pm 0.07$  Å. Deeper vertical distances have larger error bars because of electron attenuation. Other non-structural parameters are  $V_{oi} = 4.68$ eV, and isotropic thermal vibration amplitudes  $v_1 = 0.12$  Å for the topmost layer atoms and  $v_b = 0.07$  Å for subsurface atoms, determined within a Debye-Waller-like approach. These values yield a minimum  $R_{P_{min}} = 0.085 \pm 0.013$  for a database of size  $\Delta E = 1600$ eV, where the R-factor variance is  $var(R_{P,min}) = R_{P,min} \sqrt{\frac{8V_{oi}}{\Delta E}}$  [11]. If the reduced database ( $\Delta E = 940 \text{ eV}$ ) is used, the same structure yields  $R_{P,min} = 0.15 \pm 0.03$ . We shall use this value to define the lowest boundary for exiting conditions of the statistical searches as  $R_P < R_{P,min} + var(R_{P,min})$ . Thus, for N=6 interlayer distances, the search is finished when  $R_P < 0.18$ . For smaller values of N, best fit structures yield larger  $R_{P,min}$  values:  $R_{P,min}(N=2) = 0.18$ ,  $R_{P,min}(N=3) = 0.17$ ,  $R_{P,min}(N=4) = 0.16$  and  $R_{P,min}(N=5)=0.16$ . The corresponding variance is  $\sim 0.03$ . Therefore, the following search exiting  $R_P$  values are used: 0.21 for N=2, 0.20 for N=3 and 0.19 for N=4,5.

In order to examine the behaviour of SO when handling real experimental data, we perform  $N_{stat}$  equivalent statistical searches using random starting points inside the search space, that will be used to evaluate averages. The search space consists of the topmost N inter-layer spacings of the Cu(100) surface, plus a constant value of  $V_{or}$  in the range 2-10 eV. Intensities are calculated in the energy range 55-410 eV with a step of 3 eV. Thus, the nonlinear systems of equations in the inner shell contain N+1 unknowns. Size effects have been modelled using two different search spaces of hypervolume 1  $\mathring{A}^N$  and  $(0.4\ \mathring{A})^N$ . The latter corresponds to the usual Tensor LEED applicability range.

Fig. 1 shows the  $R_P$  map for N=2 in the 1 Å<sup>2</sup> space. The global minimum is located  $(d_1, d_2) = (1.75, 1.81)$  Å with  $R_P = 0.18$ , and also a few secondary deep wells with  $R_P \lesssim 0.5$  are found. The surface shows a strong corrugation, and several local shallow minima with  $R_P \sim 0.8$ . As pointed out by Rous *et al* [16], this topological feature is independent of the surface structure. Secondary deep minima are due to Bragg coincidences with the actual structure and shallow minima are due to random correlation between experimental and theoretical spectra. The  $(0.4 \text{ Å})^2$  area boundaries are 1.57 Å and 1.97 Å, respectively. This search space contains several local minima but no secondary deep minima. Due to these differences in the  $R_P$  hypersurface topography, differences in the scaling law between both search spaces appear.

The effect of different degrees of overdetermination in the inner shell has been benchmarked for N=3. Fig. 2 shows the PDF of the computing effort t at values  $N_{ov}=1,2,3$  on both search volumes.  $N_{stat}=100$  and 50 on 1 Å<sup>3</sup> and (0.4 Å)<sup>3</sup>, respectively. In both cases,  $N_{ov}=2$  appears to be a well converged value, and further

overdetermination does not cause significant improvements in the efficiency.

For  $(0.4 \text{ Å})^3$ , the PDF's are non-symmetric, show a peak at low values of t and decay exponentially. The histograms can be fitted by least-squares to Eq. 9. If plain SO is used, i.e.  $N_{ov} = 1$ , decay is clearly slower. However, in the 1 Å<sup>3</sup> volume the  $N_{ov} = 1, 3$  histograms do not show a peak, and are better fitted by an exponential function  $\lambda e^{-\lambda t}$ . Table 1 contains some average values obtained from those PDF's.

Fig. 3 shows a log-log plot of the scaling law, represented in terms of the average computing effort  $\langle t \rangle$  as a function of N. In both cases, the scaling law can be fitted to a polynomial law,  $\langle t \rangle \propto N^{\beta}$ , where  $\beta = 3.45 \pm 0.02$  for  $1 \text{Å}^N$  and  $\beta = 2.3 \pm 0.2$  for  $(0.4 \text{Å})^N$ . Averages are calculated upon  $N_{stat} = 20$  and 40 for each N value, respectively. Fig. 4 shows the average the number of  $R_P$  evaluations,  $\langle t_{R_P} \rangle$ , at different N. Least-squares fits to  $\langle t_{R_P} \rangle \propto N^{\beta_{R_P}}$  yield exponents  $\beta_{R_P} = 1.93 \pm 0.07$  for  $1 \text{ Å}^N$  and  $\beta_{R_P} = 1.0 \pm 0.2$  for  $0.4 \text{ Å}^N$ .

# 3.2. $Fe_{0.53}Al_{0.47}(100)-1\times 1$

Annealing of alloy surfaces may result in deviations of chemical composition in the surface region w.r.t. that of the bulk. Segregation in alloy surfaces is a well-known phenomenon that has been investigated in a number of binary alloys (see, for example, Ref. [30] and references therein). In the present section, we use experimental LEED I(E) data from a  $Fe_{0.53}Al_{0.47}(100)-1 \times 1$  surface to benchmark the ability of SO to fit nongeometrical data, namely impurity concentration and/or thermal vibration amplitudes. FeAl crystal structure is of CsCl type, and the stoichiometric (100) surface consists of alternating Al and Fe layers. FeAl crystals exhibit a rich phase diagram, and the surface structure strongly depends on both the bulk stoichiometry and annealing temperature. Annealing of a Fe<sub>0.53</sub>Al<sub>0.47</sub>(100) surface at 650 K results in a  $c(2 \times 2)$  reconstruction, whilst annealing at temperatures above 880 K produces a sharp  $1 \times 1$  LEED pattern. We shall focus on the latter. Experimental details on surface preparation and spectra measurements can be found elsewhere [31, 19]. Early quantitative LEED calculations on this sample determined that the (100) face is Al-terminated, yielding a satisfactory  $R_P = 0.12$  and the following interlayer distances:  $d_1 = 1.24 \pm 0.02$  Å,  $d_2 = 1.49 \pm 0.02$ Å,  $d_3 = 1.47 \pm 0.02$  Å (the lattice parameter is  $a_0 = 2.894$  Å[31]). The same analysis yields a  $\sim 20\%$  Al impurity concentration in the second layer. Since interstitial site formation is energetically unfavourable in FeAl, this result can be attributed to the strong contraction of the outermost Fe plane.

Another possible explanation is found in the coupling between chemical and vibrational parameters, which occur in LEED intensities via the Debye-Waller factor, if the average t-matrix approximation is used [32]. Blum et al [33] revisited this surface and found that the best structure fit can be achieved by either: (i) fitting the second Fe layer concentration of Al impurities,  $c_2$  and the topmost Al atomic vibration amplitude  $v_1$ , or (ii) fitting two different vibration amplitudes  $v_1$  and  $v_2$  for the first Al and second Fe layers, respectively. Considering  $V_{or}(E)$ , too, these fits yielded improved R-factors,

 $R_P = 0.091$  and  $R_P = 0.081$ , respectively. Both fits are equally favourable, as they lie within the systematic error limits of the LEED analysis.

The experimental database consists of eight beams measured at T=120 K in the energy range 40-500 eV with a step of 3 eV, producing a total data base width  $\Delta E=1880$  eV. The beam list is { (10), (11), (20), (21), (22), (30), (31), (32) }. In the I(E) evaluations,  $l_{max}=10$  provides convergence, and the imaginary part of the inner potential is kept fixed at  $V_{oi}=6$  eV. We use  $R_P \leq 0.12$  as exiting condition for SO. Independent searches have been made for N=6 parameters:  $V_{or}$ , three structural and two non-structural parameters.  $V_{or}$  lies in the range 6-14 eV. Structural parameters are the three topmost interlayer distances,  $d_1$ ,  $d_2$ ,  $d_3$ , in the ranges  $d_1=1.16-1.46$  Å,  $d_2=1.30-1.60$  Å and  $d_3=1.36-1.56$  Å. The two non-structural parameters are, according to Blum et al [33], (i)  $c_2$  and  $v_1$ , or (ii)  $v_1$  and  $v_2$ .  $c_2$  lies in the range 0-50 % and the vibration amplitudes lie in the range 0.087 - 0.15 Å. Other vibrations are kept fixed at their bulk values, namely:  $v_b(\text{Fe})=0.09$  Å and  $v_b(\text{Al})=0.12$  Å. We have performed  $N_{stat}=40$  statistical searches for each type of fit using  $N_{ov}=2$ . The corresponding computing effort PDF's are shown in Fig. 5, and average values are  $\langle t \rangle = 570 \pm 70$  for type (i) search and  $\langle t \rangle = 470 \pm 70$  for type (ii).

### 4. Discussion

The Simultaneous Optimisation (SO) algorithm has been successfully applied to parameter extraction from highly non-linear scattered intensities in ideal scenarios, namely phase shifts retrieval from single atom electron scattering and surface structure recovery from simulated LEED I(E). In these examples, SO has performed efficiently, and works well with minimal data sets in the inner shell, i.e. it uses N datapoints  $\mathcal{S}' = (E_1, \dots, E_N)$  to fit N parameters  $\vec{x} = (x_1, \dots, x_N)$ . The inner shell solves a nonlinear system of equations, a problem which is equivalent to finding the global minimum of a projected R-factor on the  $\mathcal{S}'$  data subset. The solutions provided by the system solver must be validated against the whole data base in an outer shell by evaluating the R-factor. Therefore, SO can explore a large number of meaninful configurations in the investigated hypervolume at a low computational cost. The configurations provided by the inner shell have been shown to exhibit local minima of the R-factor at the worst case (i.e. when the system solver fails to return a solution to the system because it stagnates at a local minimum of the projected R-factor). However, when dealing with real experimental databases this is not always the case. We have found that it is not guaranteed that solutions to the partial problems correspond to local minima of the R-factor, as spurious peaks in the experimental spectra may be eventually fitted. Therefore, we generalize the implementation of SO so it can be applied to realistic situations with the same efficiency or robustness as in the theoretical case. Spurious peak fitting results in high frequency corrugation of the R-factor hypersurface. By using overdetermined systems of equations, i.e. taking  $S' = (E_1, \dots, E_{N_{eq}})$  with  $N_{eq} > N$ , SO can filter out those corrugations, since the projected R-factor topography captures more

features of the global R-factor. This method has been benchmarked using experimental LEED I(E) curves in the Cu(100) surface and a search space of hypervolume 1 Å<sup>N</sup>. After determining the optimum oversampling degree, it is found that the computational effort scales polynomically as  $N^{\beta}$ , with exponent  $\beta = 3.45 \pm 0.02$ , a value which is of the same order as the previously found exponent for SO working under ideal noiseless condition. It is noteworthy that only  $N_{eq} = 2(N+1)$  are needed in the inner shell to achieve convergence in the search success rate, as shown in Fig. 2 for N=3. Since  $N_{eq} << N_D$ , it is ensured that the number of different system choices is high enough to ensure ergodicity.

As expected, the search space size affects the scaling behaviour. In particular, if we restrict the hypervolume sufficiently enough to contain only one minimum inside (local or not) and if we use gradient-directed methods, we expect an ideal scaling,  $\beta = 2$ . It is interesting to notice that the best possible scaling law,  $\beta = 1$ , implies that  $R_P$  depends linearly on the parameters, and the searching hypervolume is such that the minimum stays in one of the corners of the simplex. These are too restrictive conditions, and are not likely to happen for complex problems. On the other hand, the 1  $\mathring{A}^N$  hypervolume size is big enough to contain in most of the cases several local minima in the  $R_P$  hypersurface [20], as we have confirmed in our calculations (Fig. 1). Therefore, it is not surprising that by reducing the searching hypervolume the exponent decreases towars the ideal value of 2. Statistical search averages yield an exponent  $\beta = 2.3 \pm 0.2$ for the same system in a  $(0.4 \text{ Å})^N$  hypervolume (see Fig. 3), which is similar to the value  $\beta = 2.5$  found for RSA working in a space of the same size [19]. As shown in Fig. 4, the number of  $R_P$  evaluations in the outer shell is a linear function of N in the smaller space, while it is quadratic in 1  $Å^N$ . Therefore, we conclude that the total scaling law of SO in the former case appears to be dominated by the inner shell, explaining the high efficiency of the method.

We have also tested the applicability of SO to non-geometrical parameter retrieval. Introducing  $V_{or}$  in the search as an additional parameter does not alter significantly the scaling law in the Cu(100) example. Other non-structural parameters, such as vibrational and chemical ones, can be optimised by SO. In the studied application, Fe<sub>0.53</sub>Al<sub>0.47</sub>(100), Al intermixing in the second layer and surface vibration amplitude are satisfactorily characterised by SO using  $\langle t \rangle = 570$ , which is more than the reported value for RSA algorithm (345 trial structures) working under equivalent search conditions [19]. Interestingly, corrugation in the  $R_P$  topography is mainly due to structural parameters. Fig. 6 shows  $R_P$  as a function of chemical and/or vibrational parameters with structural parameters fixed at their best fit values. Both maps have a deep flat minimum. Therefore, we expect that statistics in this example are biased by the search of interlayer spacings.

#### 5. Conclusions

The Simultaneous Optimisation (SO) algorithm has been applied to experimental LEED data and characterized in detail so it can be compared with other alternatives in the literature. The use of noisy experimental LEED data does not imply a reduction of reliability or efficiency w.r.t. its performance on ideal noiseless databases [22]. SO brings together two apparently contradictory widespread concepts in structural work by LEED: while large parts of the database are redundant and promising structures can be obtained from a reduced number of datapoints [34], one must keep in mind that only by using large experimental databases high accuracy and reliability can be obtained. Finally, the use of the ergodic principle confers global character to the search [20].

# Acknowledgements

This work has been financed by the CYCIT (MAT-2005-3866). MBR acknowledges Spanish CSIC for a I3P fellowship. The Barcelona Supercomputer Center (http://www.bsc.es/) is acknowledged for computing time.

**Figure 1.**  $R_P$  map for Cu(100) data in a search space of area 1 Å<sup>2</sup>. Parameters  $(d_1, d_2)$  are two topmost interlayer spacings. Real part of the inner potential is also optimised for each structure. The global minimum  $(R_{P,min}=0.18)$  is located at  $(d_1, d_2)=(1.75, 1.81)$  Å. Secondary deep minima appear at  $(d_1, d_2)=(2.14, 1.45)$  Å with  $R_P=0.42$ ,  $(d_1, d_2)=(1.75, 1.48)$  Å with  $R_P=0.48$ , and  $(d_1, d_2)=(1.30, 2.20)$  Å with  $R_P=0.50$ .

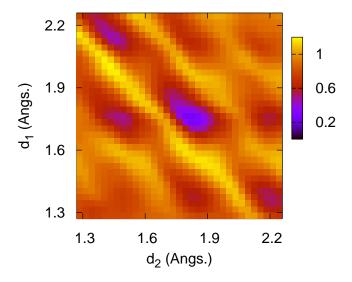
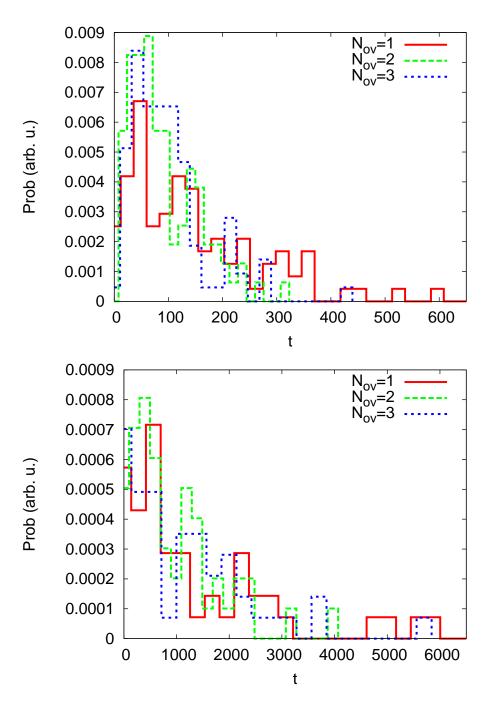
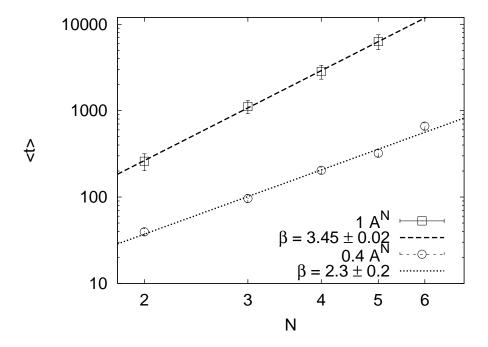


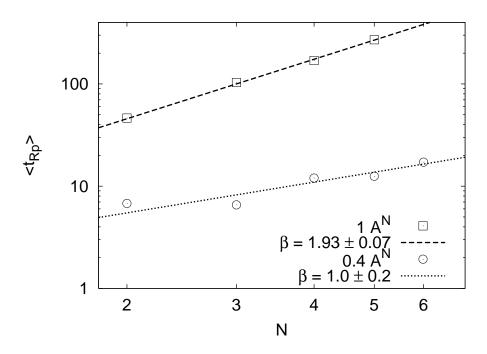
Figure 2. Panel (a): Cu(100) data in a search space of size 0.4 Å with  $N_{stat} = 100$ . Panel (b): Cu(100) data in a search space of size 1 Å with  $N_{stat} = 50$ . Search parameters are three topmost interlayer spacings. Real part of the inner potential is also optimised for each structure. Searches are made under different inner shell oversampling conditions:  $N_{ov} = 1$  (solid line), 2 (dashed) and 3 (dotted).



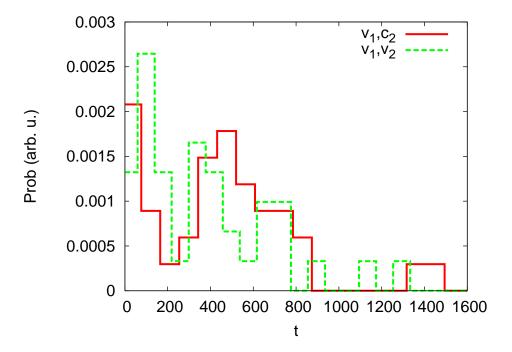
**Figure 3.** Scaling law for Cu(100) data with search spaces of sizes 1 Å<sup>N</sup> (squares) and  $(0.4 \text{ Å})^N$  (circles) together with the corresponding least-squares fits to a polynomial law.



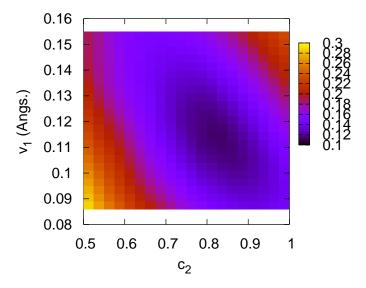
**Figure 4.** Scaling law for the number of R-factor evaluations in Cu(100) data with search spaces of sizes 1 Å<sup>N</sup> (squares) and  $(0.4 \text{ Å})^N$  (circles) together with the corresponding least-squares fits to a polynomial law.

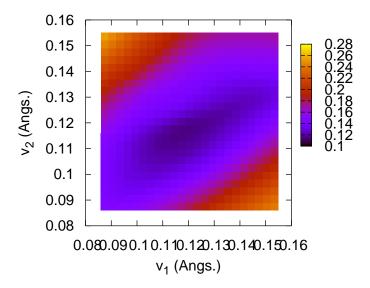


**Figure 5.** PDF's of the computing effort in FeAl searches using two types of non-structural parameters. Solid line corresponds to fitting  $c_2$  and  $v_1$ , and dotted line to fitting  $v_1$  and  $v_2$ . Both histograms are built upon  $N_{stat} = 40$  independent searches.



**Figure 6.**  $R_P$  map for Fe<sub>0.53</sub>Al<sub>0.47</sub>(100) data as a function of (a)  $(v_1, c_2)$  or (b)  $(v_1, v_2)$  (see text), with structural parameters fixed at their best fit values.





**Table 1.** Average t values corresponding to histograms of Fig. 1.  $\langle t \rangle$  is the histogram average,  $\langle t \rangle_E = 2/w$  is the expected value obtained after a least-squares fit of the histograms to a function  $p(t) = w^2 t e^{-wt}$  and  $\langle t \rangle_P$  (P = 0.90, 0.9998) are the average t values needed to get a successful search with P probability. The values marked by (\*) have been obtained from a fit to  $p(t) = \lambda e^{-\lambda t}$ .

		$0.4~\mathrm{\AA^3}$				$1 \text{ Å}^3$		
$N_{ov}$	$\langle t \rangle$	$\langle t \rangle_E$	$\langle t \rangle_{0.90}$	$\langle t \rangle_{0.9998}$	$\langle t  angle$	$\langle t \rangle_E$	$\langle t \rangle_{0.90}$	$\langle t \rangle_{0.9998}$
1	$160 \pm 10$	$146 \pm 6$	284	805	$1620 \pm 150 \ (*)$	$1270 \pm 160$	3729 (*)	13793 (*)
2	$100 \pm 6$	$88 \pm 4$	170	482	$1050 \pm 120$	$960 \pm 100$	1873	5297
3	$110 \pm 7$	$97 \pm 5$	188	532	$1550 \pm 150 \ (*)$	$1600 \pm 300$	3578 (*)	13235 (*)

- [1] J B Pendry. Low-Energy Electron Diffraction. Academic, London, 1974.
- [2] M A Van Hove, V H Weinberg, and C-M Chan. Low-Energy Electron Diffraction. Springer, Berlin, 1986.
- [3] M A Van Hove and S Y Tong. Surface crystallography by LEED. Springer, Berlin, 1979.
- [4] D P Woodruff and A M Bradshaw. Adsorbate structure determination on surfaces using photoelectron diffraction. *Rep. Prog. Phys.*, 57:1029, 1994.
- [5] G H Stout and L H Jensen. X-ray structure determination: a practical guide. John Wiley and sons, New York, 1989.
- [6] I. K. Robinson and D. J. Tweet. Surface x-ray diffraction. Rep. Prog. Phys, 55:599, 1992.
- [7] D. Norman. X-ray absorption spectroscopy (exafs and xanes) at surfaces. J. Phys. C: Solid St. Phys., 19:3273, 1986.
- [8] D P Woodruff and T A Delchar. Modern Techniques of Surface Science. Cambridge University Press, Cambridge, 1994.
- [9] A Zangwill. Physics at Surfaces. Cambridge University Press, Cambridge, 1998.
- [10] C G Kinniburgh and J B Pendry. The phase problem in leed. J. Phys. C: Solid St. Phys., 11:2415, 1978.
- [11] J B Pendry. Reliability factors for leed calculations. J. Phys. C: Solid St. Phys., 13:937, 1980.
- [12] P J Rous, J B Pendry, D K Saldin, K Heinz, K Müller, and N Bickel. Tensor leed: A technique for high-speed surface-structure determination. *Phys. Rev. Lett.*, 57:2951, 1986.
- [13] J B Pendry, K Heinz, and W Oed. Direct methods in surface crystallography. *Phys. Rev. Lett.*, 61:2953, 1988.
- [14] D K Saldin, X Chen, J A Vamvakas, M Ott, H Wedler, K Reuter, K Heinz, and P L de Andres. Holographic leed: A review of recent progress. Surf. Rev. and Lett., 4:991, 1997.
- [15] P G Cowell and V E de Carvalho. Unconstrained optimization in surface crystallography by leed preliminary results of its application to cdte(110). Surf. Sci., 187:175–193, 1987.
- [16] P J Rous, M A Van Hove, and G A Somorjai. Directed search methods for surface structure determination by leed. *Surf. Sci.*, 226:15, 1990.
- [17] D L Adams. A simple and effective procedure for the refinement of surface structure in leed. Surf. Sci., 519:157–172, 2002.
- [18] R Döll and M A Van Hove. Global optimization in leed structure determination using genetic algorithms. Surf. Sci., 355:L393, 1996.
- [19] M Kottcke and K Heinz. A new approach to automated structure optimization in leed intensity analysis. Surf. Sci., 376:352, 1997.
- [20] P J Rous. A global approach to the search problem in surface crystallography by low-energy electron diffraction. *Surf. Sci.*, 296:358, 1993.
- [21] V B Nascimento, V E de Carvalho, and et al. The fast simulated annealing algorithm applied to the search problem in leed. Surf. Sci., 487:15, 2001.
- [22] M Blanco-Rey and P L de Andres. Surface diffraction structure determination from combinatorial simultaneous optimization. Surf. Sci., 600:L91, 2006.
- [23] C G Broyden. A class of methods for solving non-linear simultaneous equations. *Mathematics of Computation*, 19:577, 1965.
- [24] G Kleinle, W Moritz, D L Adams, and G Ertl. A novel procedure for fast surface structural-analysis based on leed intensity data. Surf. Sci., 219:L637–L645, 1989.
- [25] R. H. J. M. Otten and L. P. P. P. van Ginneken. The Annealing Algorithm. Kluwer, Boston, 1989.
- [26] W H Press, S A Teukolsky, and W T Vetterling. Numerical Recipes. Cambridge University Press, Cambridge, 2002.
- [27] S Müller, A Kinne, M Kottcke, R Metzler, P Bayer, L Hammer, and K Heinz. In-plane lattice reconstruction of cu(100). *Phys. Rev. Lett.*, 75:2859, 1995.
- [28] S Walter, V Blum, L Hammer, S Müller, K Heinz, and M Giesen. The role of an energy-dependent inner potential in quantitative low-energy electron diffraction. *Surf. Sci.*, 458:155, 2000.
- [29] J Rundgren. Electron inelastic mean free path, electron attenuation length, and low-energy

- electron-diffraction theory. Phys. Rev. B, 59:5106, 1999.
- [30] L Hammer, W Meier, V Blum, and K Heinz. Equilibration processes in surfaces of the binary alloy fe-al. J. Phys.: Condens. Matter, 14:4145–4164, 2002.
- [31] M Kottcke, H Graupner, D M Zehner, L Hammer, and K Heinz. Segregation-induced subsurface restructuring of feal(100). *Phys. Rev. B*, 54:R5275, 1996.
- [32] Y Gauthier, Y Joly, R Baudoing, and J Rundgren. Surface-sandwich segregation on nondilute bimetallic alloys:  $Pt_{50}ni_{50}$  and  $pt_{78}ni_{22}$  probed by low-energy electron diffraction. *Phys. Rev. B*, 31:6216, 1985.
- [33] V Blum, L Hammer, W Meier, and K Heinz. Quantification of substitutional disorder and atomic vibrations by leed the role of parameter correlations. *Surf. Sci.*, 488:219–232, 2001.
- [34] G Kleinle, W Moritz, and G Ertl. An efficient method for leed crystallography. Surf. Sci., 238:119, 1990.